





In 3.x/4.0 the Session manager is configurable at the Application server level. So all the applications share the same SessionManager settings. This doesn't allow isolation at the application level. Different applications might have different level of requirement on session manager, like persistence, cookies etc. This will be useful for the ASPs who run more than one application on a server. This also reduces the load the database when used with session persistence because applications do not have to share the same database for session persistence. By Default, SessionManager settings defined at the application server level are inherited by the Enterprise applications and by the web modules within them.



The Servlet 2.2 Specification limited the scope of sessions to the web module, and Servlet 2.3 does not remove this restriction. However, WebSphere 5.0 offers an extension to the specification, such that you can configure session scoping on the Enterprise application. This feature is not in the Admin Console, but available in the Application Assembly tool and in Application Developer.

There is also a configuration option to handle session persistence at the Enterprise application level and at the Web Module level. If you chose to scope sessions at the Enterprise application level, you will not be able to alter the session persistence at the module level.



- Cluster Member order based session affinity: When a clone member is marked dead, the Version 4.0 plug-in directs the request to one of the available clones randomly. In this case, there is a chance that two concurrent requests for a session after clone went down might end up in different clones. To avoid this, the Version 5.0 plug-in will maintain the cluster member list in order and rather than picking randomly. The plug-in will pick the cluster member next in its list to avoid the breaking of session affinity. This is also useful in case of URL-rewriting where cluster member information doesn't get updated if rewritten URLs aren't updated.
- Speaking of URL Rewriting, most Wireless Application Protocol (WAP) devices do not support cookies. The preferred way to track sessions for WAP devices is to use URL rewriting. However on most WAP devices, the maximum allowed URL length is 128 characters. With URL rewriting, a session identifier is added to the URL itself, effectively decreasing the space available for the actual URL and the number of parameters that can be sent on a request.
- As of Fix Pack 1, you can reduce the length of session identifier, by configuring the session ID length and clone ID on an application server's web container.



Serialize the access to the Session: As per servlet API, session can be accessed concurrently and it would be responsibility of the developer to serialize the access to the session. As a quality of service, we provide a serialized access to the session in a given JVM so that access to the thread safe attributes in session can be avoided without having to write any special code by the developer. The wait time is configurable.



We still have the capability of persisting sessions to a data table, as we did in 4.0. We also have the capability of persisting sessions to a memory location.

Memory to Memory replication can use one of three basic topologies - Single replica, client/server and peer-to-peer. Each has advantages which we examine in the slides to come.

Memory to memory session replication has all the features that database persistence has (except the database-specific options, like row size).



Configuration is very straightforward; you configure the topology, the system handles the rest.

When to persist is also configurable - at the end of each request, at a timer interval, or when directed by the application code - which is to say, when the synch() method is called.

Memory to Memory Replication traffic can be encrypted. It is available in WebSphere Network Deployment and above.

Session persistence to memory has some performance advantages over database persistence, and scales well.



WebSphere Internal Messaging is built on a small, fast publish/subscribe engine which uses a JMS-like function to communicate session information from place to place. There is a separate thread spawned for each queue. By default, the topology is peer-to-peer, but that particular topology, while requiring the least configuration, does not scale as well as the others.

One feature to permit better scaling is Partitioning, which we'll detail in a few slides.

Session data is not all that can be replicated - Dynamic Data caches can also be moved from machine to machine using WebSphere Internal Messaging.



WebSphere Internal Messaging can be configured as N-way Peer-to-Peer, as on the left. While this is the simplest to configure, there are performance implications when scaling up to large environments - because every server has to persist the sessions of every other server.

A special case of Peer-to-peer is the Single Replica option, where session info is persisted to only one other application server. This scales well, but introduces something like a single point of failure - i.e., if BOTH application servers are taken out at once (for instance, the whole node goes down), then persistence breaks.

So there is another topology option - configuring Client/Server setup, where one app server is dedicated to persist all the sessions for all the app servers. This provides failover, but then the server becomes a single point of failure. So we can set up to have two app servers at different locations both persisting the sessions for all the other app servers, as illustrated on the right.

		IBM
WebSphere Interna	al Messaging - Sec	urity
<ul> <li>Environment-&gt; Internal Replication</li> </ul>	<b>Configuration</b> RegenerateKey	
Domains->	General Properties	
<name></name>	Name	* MyCluster
Messages can be	Request Timeout	* 5
encrypted - DES	Encryption Type	* NONE
or Triple DES	DRS partition size	* NONE DES
	Single Replica	
11 HTTP Session Management		© 2002, 2003 IBM Corporation

The messaging traffic that carries session information from replicator to replicator can be encrypted for added protection in the runtime environment. A button is provided to regenerate the encryption key that is used.



Partitioning is for a large environment, so that members of one cluster need not deal with the 'chatter' from other clusters. Each cluster member persists its session data to the configured listeners - that is, all the other app servers listening to that partition ID.

This can be configured to match the environment. If it is needful that an application server be dedicated to the task of session persistence, and serve no end-user requests, it can still be set up to handle only certain partition IDs. Or each application server can listen to all the available partition IDs in a peer-to-peer scenario.

If a cluster is constructed such that not all members listen to the same partition IDs, internal logic will add the partition ID necessary to handle sessions should an application server fail.

It is suggested that each segment of a replication environment consist of at least ten partition IDs, so that cluster members have multiple paths to communicate with the application server that is persisting its session data.

These partitions/groups correspond to Topics. There is one and only one thread associated with a given subscriber for a Topic. Through testing, we selected 10 subscriber threads to handle the traffic.

Name	Description	Value	Select
Created Sessions	Number of sessions that were created	1	~
Invalidated Sessions	Number of sessions that were invalidated	1	~
Session Life Time (ms)	Average session life time in milliseconds (tim	2,122,512	~
Active Sessions	Total number of sessions that are currently be	0.0 (Avg: 0.0	
Live Sessions	Total number of sessions that are currently live	0.0 (Avg: 0.2	
No Room For New Se	Applies only to session in memory with Allow	0	
Cache Discards	Number of session objects that have been for	0	
External Read Time	Time (milliseconds) taken in reading the sess	0	
External Read Size	Size of session data read from persistent stor	0	
External Write Time	Time (milliseconds) taken in writing the sessi	0	
External Write Size	Size of session data written to persistent store	0	
Affinity Breaks	Number of http session affinities are broken, n	0	
Time Since Last Activ	The time difference in milliseconds of the prev	10,789.762	
Invalidated Via Timeout	The number of sessions that are invlidated via	1	
Activate Non-exist Ses	Number of requests for a session that no long	0	
Activate Non-exist Ses	Number of requests for a session that no long	0	

Tivoli Performance Viewer is the updated and re-branded Resource Analyzer that we used in version 4.0.

The first four items in the list were available in WebSphere 3.5.5; the remaining items are new in 5.0.

TRM







The cookie domain should match servlet URI mapping. For example, if the cookie domain is set as myCompany.com, the servlet should be accessed with that domain name. (http://mySystem.myCompany.com/myapp/servlet/sessionservlet)











Multiple instances of browsers started as sub process share the cookies, that's why session gets shared between. It is the browser behavior that is defining this.









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