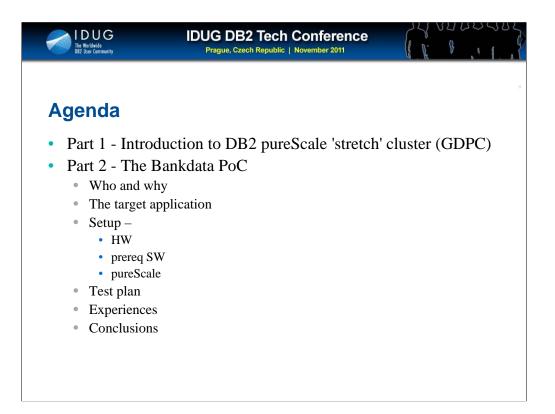
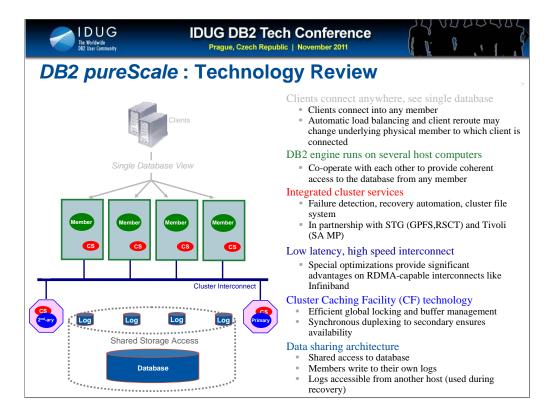
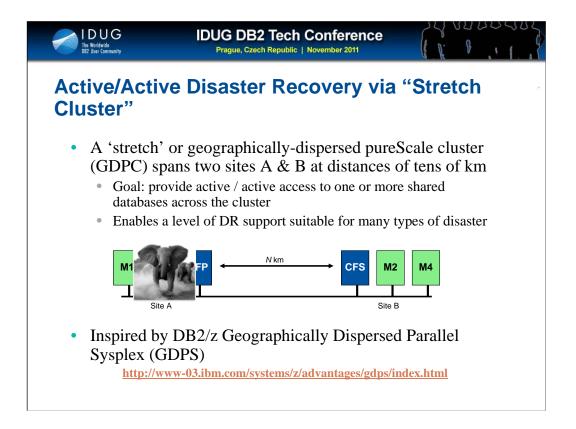


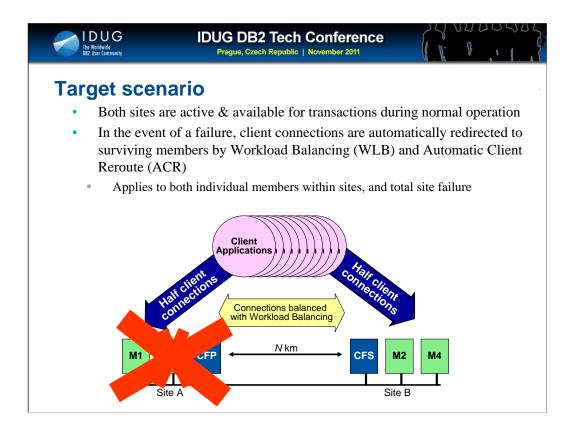
This presentation covers DB2 pureScale geographically dispersed clusters. pureScale is DB2 LUW's answer to DataSharing in DB2 z/OS to provide DB2 LUW with unlimited scalability and high availability. Until now, pureScale clusters have only been supported in configurations where the physical boxes were placed within a very limited distance of one another. By exploiting some very advanced network technology, pureScale clusters can now be 'stretched' so that they can be used in an installation where 2 sites are placed kilometers apart. In this way pureScale can take part in an "active/active" disaster recovery (DR) setup, where one surviving site can take over the workload should the other site fail, and in this way ensure maximum availability. In 2011 Bankdata/JN Data and IBM started a "Proof of Concept" to bring this setup into the Bankdata/JN Data installation. This presentation will be covering both the physical (hardware) setup, the software setup and information about all the tests of the error scenarios that were performed at Bankdata/JN Data to verify the solution.



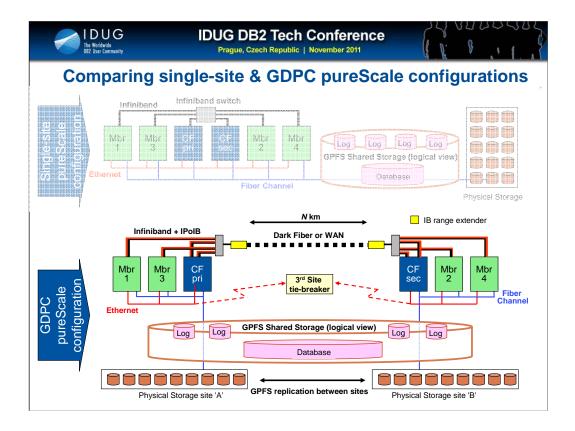




The 'geographically dispersed pureScale cluster' allows fully synchronized read/write activity concurrently at both sites. If one site should happen to go down (say, due to charging elephants, or other more prosaic problems), the other half of the cluster remains functional to continue work.



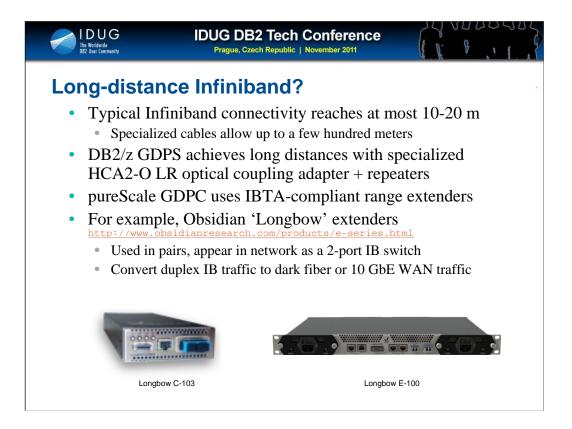
Different levels of failures – from single or multiple members, members and CF, or even an entire site – can be handled by the GDPC configuration.



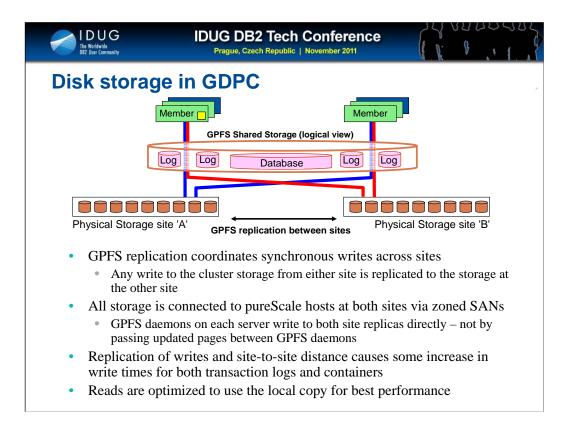
In simplest terms, GDPC basically splits a regular pureScale cluster, putting half of the compute resource at each site. Communication is maintained with advanced Infiniband extender technology, and synchronized versions of the on-disk data are maintained transparently using GPFS synchronous replication.

The 3<sup>rd</sup> site tie-breaker is required to avoid 'split brain' cases, where the network between sites might go down, and neither side can legitimately claim to be "THE" cluster afterward (or worse, they both do.) The tie-breaker is very modest – it does not have to have access to the SAN or IB. The tie-breaker *could* be located at either of the two main sites – but then there are challenges if both site A (for example) and the tiebreaker go down. Extra steps will be required to bring up the remaining site, since quorum can't be achieved with both A and tiebreaker down.

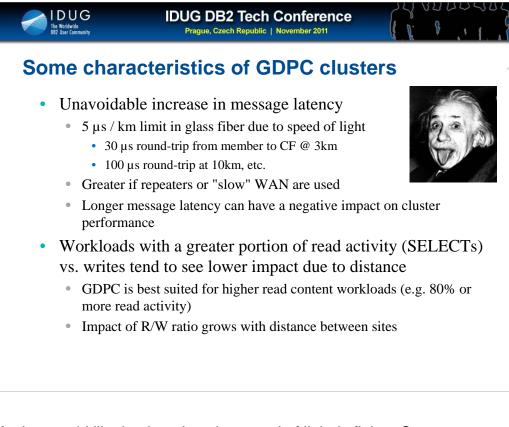
For details on GPFS synchronous replication, see GPFS Admin doc SC23-5182-02



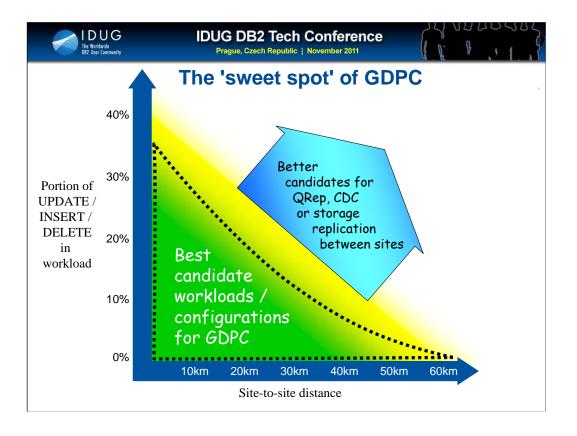
Long-distance Infiniband exists in the DB2 Sysplex world already, providing support for mainframe-based distributed clusters (GDPS). pureScale utilizes IBTA-standard Infiniband adapters and switches, which can be extended over long distances with devices such as the Obsidian Longbow IB extenders.



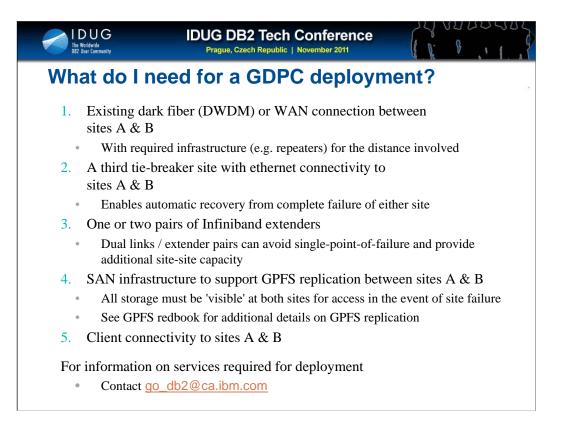
DR isn't DR unless each site can continue operation without the other. And that means that two copies of the database must exist – one at each site. GPFS synchronous replication keeps them in sync, but a zoned SAN setup provides the underlying infrastructure for the systems at each site to access the disks at both sites.



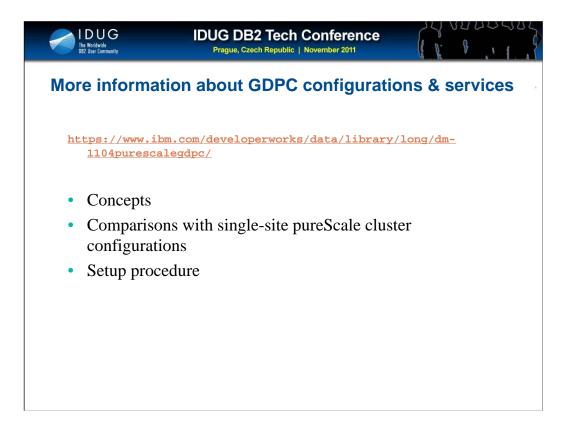
Much as we'd like it otherwise, the speed of light is finite. So once we start adding distance between sites, messaging delays due to the speed of light start to creep in. These can be quite insigificant at very short distances, however they can add up as the sites get to be 20 or 30 or more km apart. Because database write operations require more message traffic on average than read operations, the nature of the workload (read heavy, or write heavy?) is an additional factor in maximum practical distance between sites.



pureScale supports GDPC as well as other DR solutions, such as Q Replication (QRep) and Change Data Capture (CDC). The 'sweet spot' of GDPC, where it's the most suitable choice, typically involves relatively close site-to-site distances, and higher read ratios.

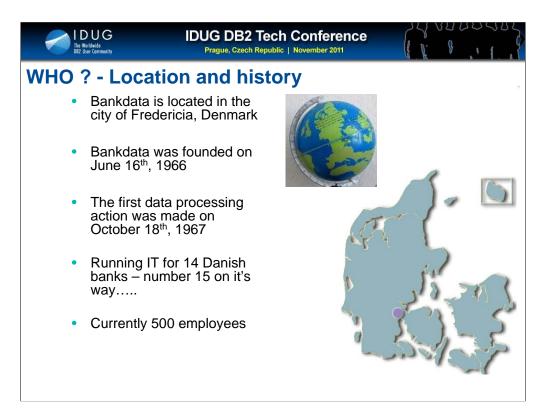


GDPC is fundamentally a 'normal' pureScale cluster stretched over two sites – so the core system requirements of hardware and software are very similar. The main differences on top of this include (1) a highbandwidth, low-latency WAN or dark fiber connection between sites, (2) Infiniband extenders to span that distance, (3) zoned SAN storage to provide disaster toleration between sites, and (4) GPFS replication to keep storage content in sync.

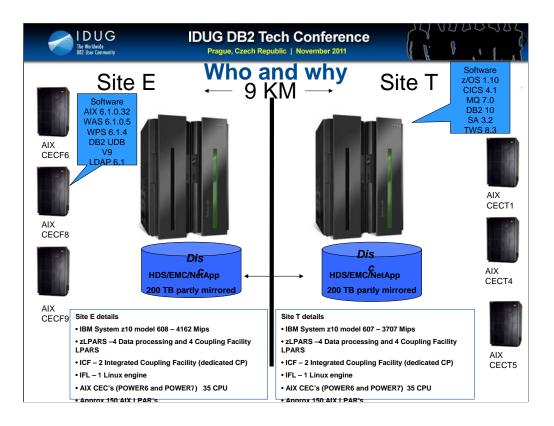


The referenced whitepaper goes into detail on how a GDPC is configured, how it's different from a regular pureScale cluster, etc. Definitely required reading for anyone interested in a GDPC deployment!

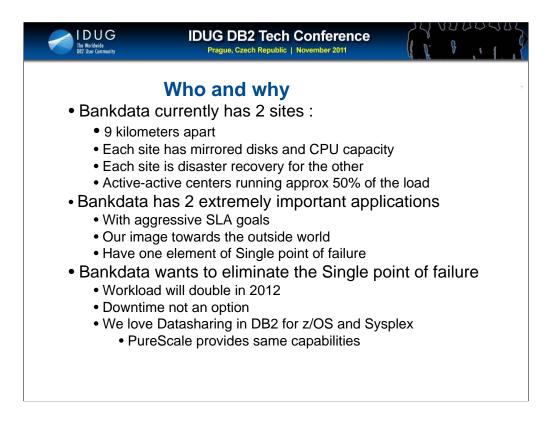




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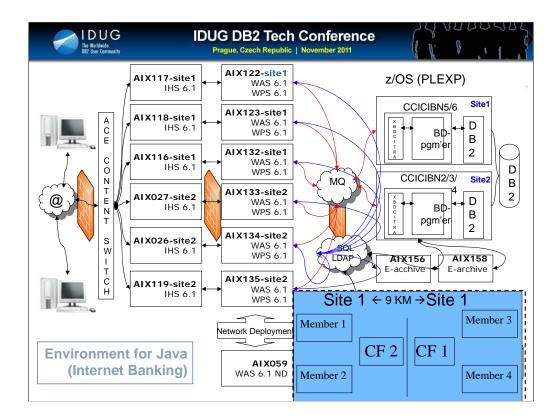
Hardware and computer center physical layout.



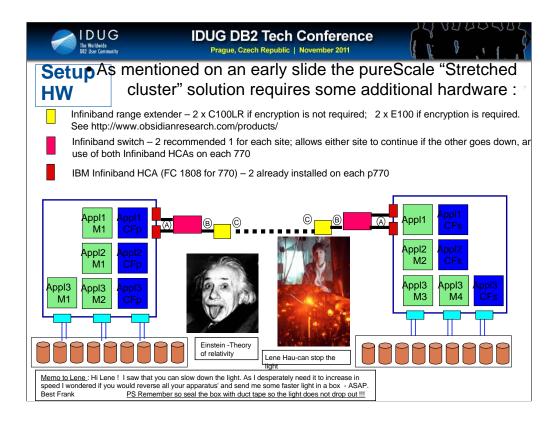
Bankdata has an active-active politics concerning computer centers. This means that for a "level-1" disaster (if there is such a thing) where we "just" lose one center processing should be able to continue unaffected. On z/OS this is pretty easy as we have SYSPLEX and DB2 DataSharing etc so by placing sufficient mainframe capacity and mirroring all disks this is pretty straight forward.

However on Windows and partly on AIX this is more difficult. Our main customer applications is running with most of the data on DB2 on z/OS using CICS transactions as ""WEB service calls"". Presentation layer and transaction driver is WAS/WPS running Java code on AIX. We have 50% of the WAS/WPS instances on either computer center but there is a small but highly active part of the data in DB2 UDB on AIX. Until now this has been a SPoF as we have had to make this an active/inactive DB2 UDB solution ("HACMP-like"). So if we loose one center we will have to wait for the DB2 UDB to be activated on the other side.

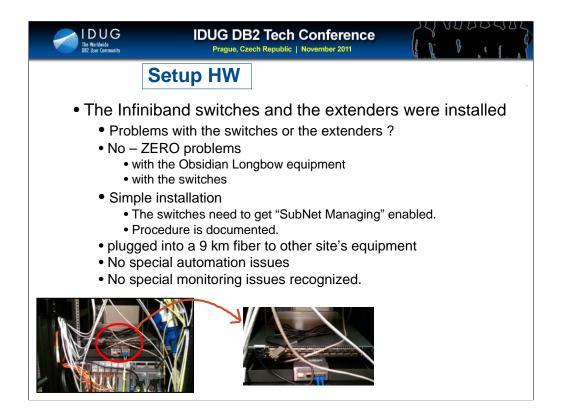
Similar upgrades and maintenance on the DB2 AIX will give a outage, though planned. We have in many years looked at DataSharing in DB2 on z/OS with envy when we are wearing our DB2 UDB glasses !!!



Just a picture on one of the main appliaction to highlight where the problem with SPoF is !!



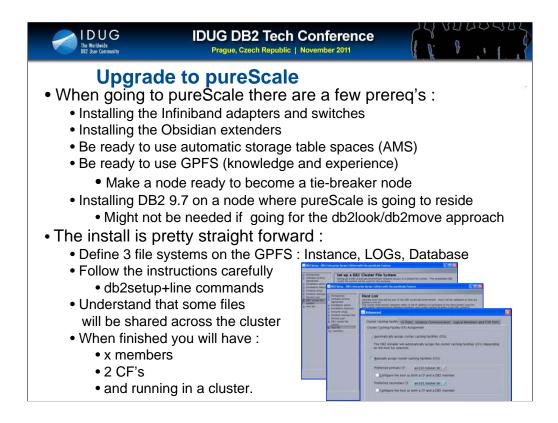
As understood by most people now pureScale brings DB2 on z/OS data sharing into DB2 UDB. Not implemented in the hardware but in a approximate version simulated in software components. Until now the two members of the cluster had to be 'within a data center' but with the Stretched Cluster solution WITH the usage of additional hardware we can now extent the distance – penalty being the speed of light .....



The Obsidian Longbow equipment is beautiful ! I am fond of beautifully solutions and I felt quite sad when we plugged all the fibers into them and stacked them in these messy racks ! Throughout the POC we observed no problems with all the equipment. We did no special monitoring because of this but we think that it is in fact possible to extract figures from the switches.

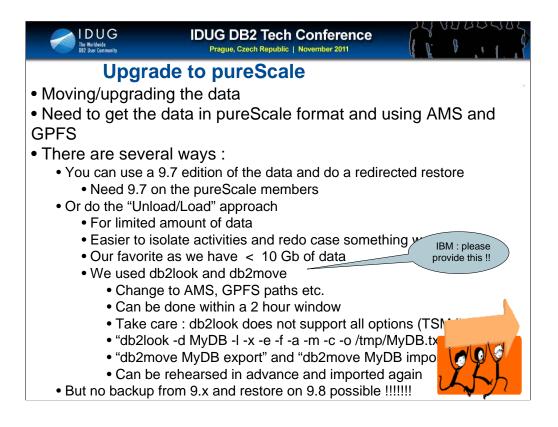
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| • Hov | v does these                             | new versions       | fit into the p        | ureScale              |      |
| requi | Product                                  | Current<br>version | Planned<br>upgrade to | pureScale<br>supports |      |
|       | WAS                                      | 6.1                | 7                     | 6.1 >                 |      |
|       | WPS                                      | 6.1                | 7                     | 7 =>                  |      |
|       | DB2 (AIX)                                | 9.1                | 9.7                   | (9.8)                 |      |
|       | LDAP (TDS)                               | 6.1.5              | 6.2                   | Next Release ?        |      |
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There are some requirement for the software supported by pureScale. In our environment we needed to upgrade WAS/WPS anyway so it did not present a real problem. The closest we came to a problem was TDS or LDAP where we could not get the required version. There are however several ways to get passed this.....



You have to realize that if you are a DB2 for z/OS datasharing freak things are a bit more complicated as the solution builds upon and around other building bricks and technology that you have to consider to which extent you have to master these....

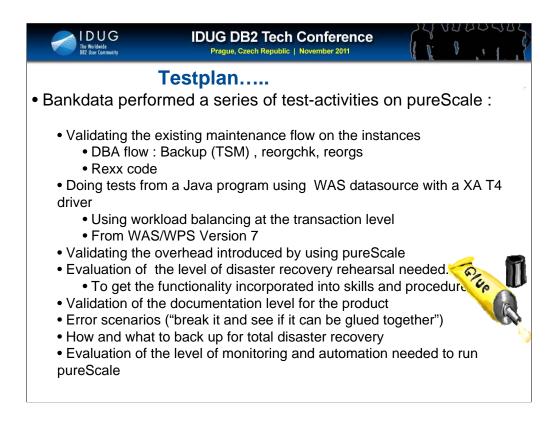
For instance GPFS and if you have never user Automated Storage Managed Table Spaces. The thing is that the setup of much of this is done once and build as much into the pureScale installation procedure as possible but you still have to understand that in a disaster situation you have to be able to operate all these technology components.



Now we need some data in our newly defined pureScale instance. The filesystems has to be GPFS and the table spaces AMS. If you will use the redirected restore you will need a DB2 9.7 on the system. Another way is to unload the data from the 9.x database and use a "load" on the pureScale instance. There are some clear advantages in this, for instance can you build the pureScale environment way before and do several test-loads before the final day arrives....

To do the unload/load a db2look and db2move is very elegant. You will however need to do some manual definitions at the pureScale side and understand that not all will be moved by db2look....

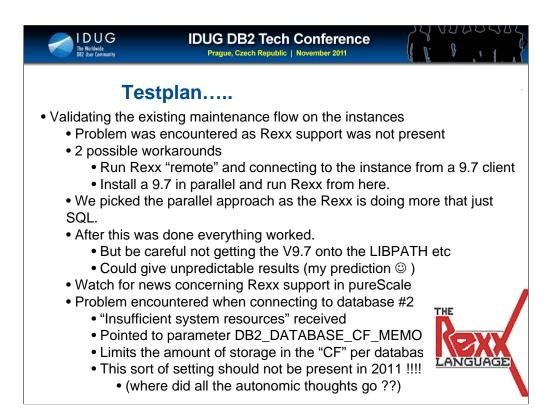
But it can be tested and rehearsed and build into a script. It would have been elegant if it was build into the pureScale installation as a script so every customer did not have to find the same pitfalls.... But it will soon be Christmas...



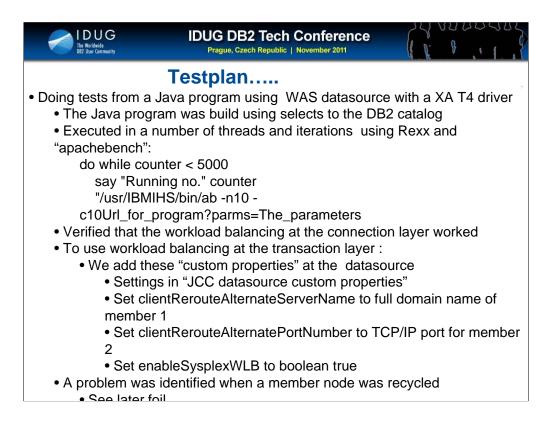
Okay, we have installed all the components and moved the data. Time is now right for testing the real reason why we do this : will pureScale let the workload continue with acceptable slowdown or interrupt without any manual intervention ? Will we see problems in the existing 'housekeeping' flow ? Will we impose an overhead on every transaction ? Is the doc understandable and sufficient ? Do we need additional monitoring ?

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And most important will the systen survive all these more or less expected scenarios where we today will have an outage ??



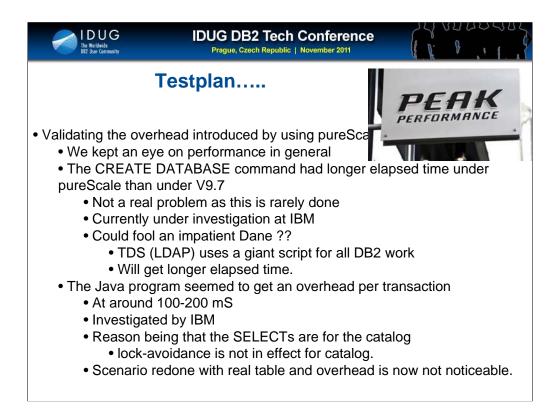
Rexx code gave us some problems as this was not included in the version of pureScale that we tested. We expect it to be included in the near future but until them we will need a filesystem with the V9.7 code to allow Rexx. But do take care not to get V9.7 code into play in the pureScale instance !!!



When conducting test cases where you have to compare performance and verify specific scenarios you have 2 different routes to go. You can execute a 'real' workload or you can write some test programs that will fire of the canned set of workload. The later one will normally be the best way to go as it will allow you to redo the same scenario as often as you will to validate different settings and fix-levels. We made a small test driver simply doing some catalog lookup and executes this under a type 4 XA driver using the apache bench. This allow you to start a number of parallel threads executing the same application and by wrapping this in a script (Rexx or shell) you can also control how many loops you will conduct. We used this to verify the stability, the performance and the workload balancing in the product.

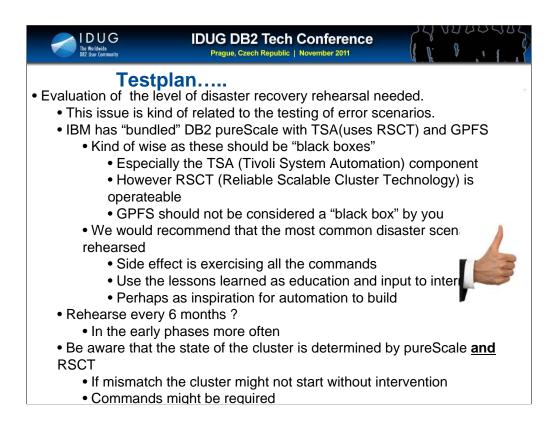
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To set the workload balancing at the transaction layer you need to set some properties in the WAS admin console under "custom properties".



We were surprised that a common thing like a CREATE DATABASE took much longer time under pureScale than under DB2 UDB 9.7. As this is very rarely done the biggest danger in this is that users (installers) get impatient and cancel the process.

What worried us more was that we saw a big overhead on all transactions. IBM investigated and it proofed to be a consequence of the Java program selecting in the catalog where pureScale's lock-avoidance is not in effect. So properly/surely not a problem in a real world.....

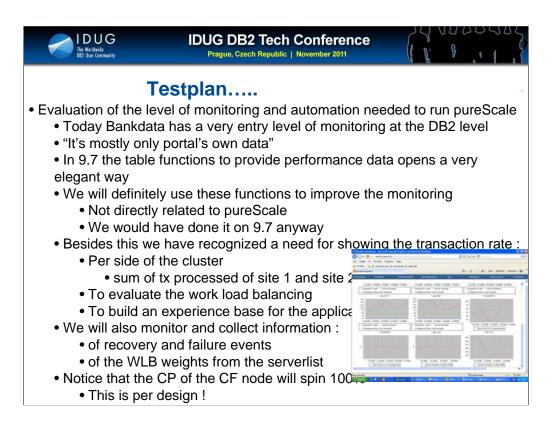


During the testing of all the possible and impossible error scenarios we saw that our knowledge of the underlying components GPFS, RSCT and TSA was not good enough as the health of the system involves decisions and actions in all these components. We advise new users of pureScale to invest time in getting familiar with these components and their commands. Understand that this is an investment because you might newer get in to a situation where this knowledge will be needed. But in a disaster knowledge can make the difference....

So please rehearse all the unthinkable scenarios in time of peace !!!

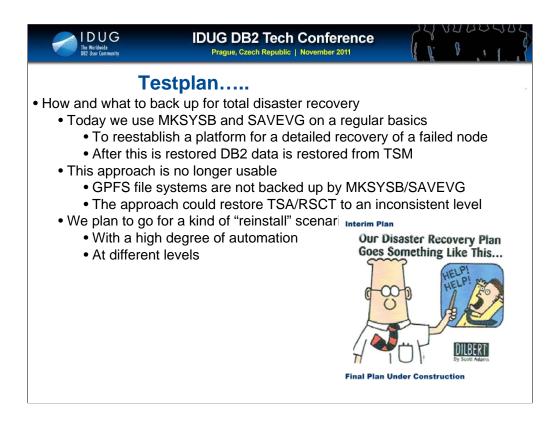


For a systems programmer this is the worst task of all : reading documentation. We went through much of it and our feeling was that the base doc was well written and adequate but there was surely a need for more technical doc and white papers to help in understanding the internals in all the related components and how these interact with pureScale. This has been taken "ad notam" by IBM....



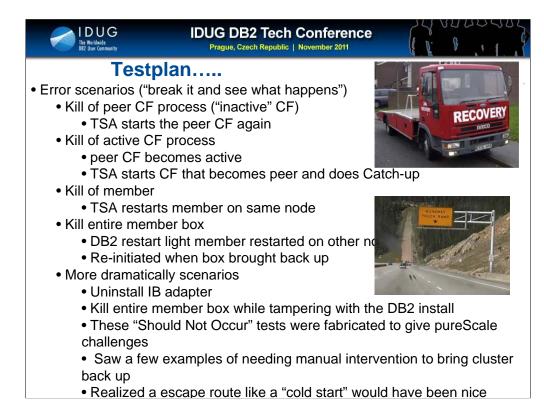
When looking at our feelings concerning monitoring it is important to understand that we were running a DB2 UDB 9.1.x at that time. In 9.7 all these wonderful table functions arrived that will give you the power of seeing all the performance metrics in an easy way. Our monitoring today just graphs the transaction rates and some thread info with alerts case something goes wrong. If we decided NOT to go to pureScale we would go to DB2 UDB 9.7 and we would change to using all these metrics to modernize our monitoring of the DB2 UDB.

However if we look at our monitoring for our CICS on z/OS running SYSPLEX we simply show one plot with the sum of executed transactions on ALL the CICS regions running on each "side". If a CICS region crashes it will not be reflected as long as the surviving members are capable of processing the incoming transactions and System Automation will simply restart the failing CICS. With pureScale we will do something similar : plot the database activity on a plot per "side" (read : computer center) and we will ensure to collect available information of events as member down and recovery events. We also expect to show information of the weights used by workload balancing.



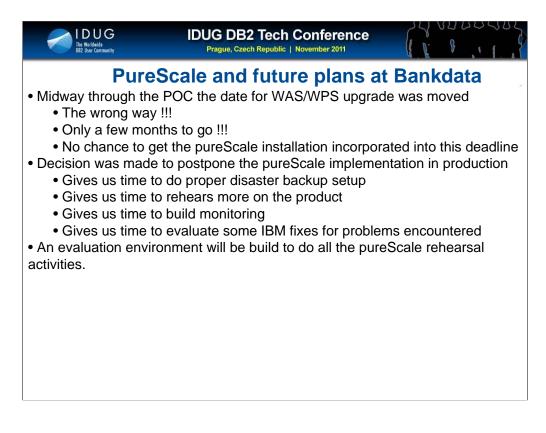
For the planning for the total disaster (loss of both centers) we today use a MKSYSB and SAVEVG to be able to rebuild the nodes in the applications.

However going to GPFS this is no longer a valid approach and the most likely strategy is to have some sort of reinstall plans and automation. This is one of the areas where a white paper will be very welcome....

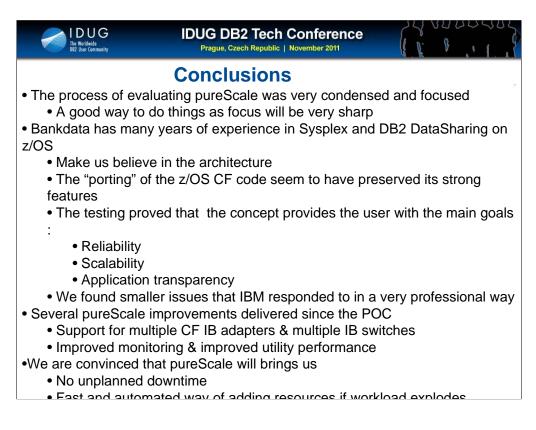


The final heat : testing all the scenarios that is thinkable and unthinkable. We tested all the thinkable ones and these did exactly what they should with no or with smaller interruption in service.

We also did things that will never happen and some of these indeed could bring the system in a state where it could not start automatically. IBM has been very aggressive in finding ways to handle these situations as well inside the code even though is was "ShouldNotOccur" situations.



Where are we today? In the Bankdata plans for upgrading these 2 vital application there were included enough slack to do a migration to pureScale. However these plans were changed so the applications needed to be upgraded earlier that expected. We therefore had to postpone the migration to pureScale and are at the moment using this extra time to get even more familiar with the product and verify some of the improvements to come.



The conclusion is that the pureScale product is working excellent and that it is working very similar to DataSharing on z/OS as it is to some degree a porting of the code and ideas from this to the DB2 UDB.

We met a very dedicated and skilled staff at the IBM Toronto Lab that responded quickly and competent on our problems and silly questions.



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