## Pascal Users Group

# Pascal News 

Communications about the Programming Language Pascal by Pascalers

- APL Scanner
- Computer Generated Population Pyramids
- Path Pascal
- Introduction to Modula-2
- Validation Suite Reports
- Announcements

- Pascal News is the official but informal publication of the User's Group.

Purpose: The Pascal User's Group (PUG) promotes the use of the programming language Pascal as well as the ideas behind Pascal through the vehicle of Pascal News. PUG is intentionally designed to be non political, and as such, it is not an "entity" which takes stands on issues or support causes or other efforts however well-intentioned. Informality is our guiding principle; there are no officers or meetings of PUG.

The increasing availability of Pascal makes it a viable alternative for software production and justifies its further use. We all strive to make using Pascal a respectable activity.

Membership: Anyone can join PUG, particularly the Pascal user, teacher, maintainer, implementor, distributor, or just plain fan. Memberships from libraries are also encouraged. See the COUPON for details.

- Pascal News is produced 4 times during a year; January, April, July October.
- ALL THE NEWS THAT'S FIT, WE PRINT. Please send material (brevity is a virtue) for Pascal News singlespaced and camera-ready (use dark ribbon and 15.5 cm lines!)
- Remember: ALL LETTERS TO US WILL BE PRINTED UNLESS THEY CONTAIN A REQUEST TO THE CONTRARY.
- Pascal News is divided into flexible sections:

POLICY - explains the way we do things (ALL-PURPOSE COUPON, etc.)
EDITOR'S CONTRIBUTION - passes along the opinion and point of view of the editor together with changes in the mechanics of PUG operation, etc.

APPLICATIONS - presents and documents source programs written in Pascal for various algorithms, and software tools for a Pascal environment; news of significant applications programs. Also critiques regarding program/algorithm certification, performance, standards conformance, style, output convenience, and general design.

ARTICLES - contains formal, submitted contributions (such as Pascal philosophy, use of Pascal as a teaching tool, use of Pascal at different computer installations, how to promote Pascal, etc.).

OPEN FORUM FOR MEMBERS - contains short, informal correspondence among members which is of interest to the readership of Pascal News.

IMPLEMENTATION NOTES - reports news of Pascal implementations: contacts for maintainers, implementors, distributors, and documentors of various implementations as well as where to send bug reports. Qualitative and quantitative descriptions and comparisons of various implementations are publicized. Sections contain information about Portable Pascals, Pascal Variants, Feature-Implementation Notes, and Machine-Dependent Implementations.

VALIDATION SUITE REPORTS - reports performance of various compilers against standard Pascal ISO 7185.

# Pascal News 

Communications about the Programming Language Pascal by Pascalers

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Charles Gaffney Publisher and Editor
The Pascal Newsletter is published by the Pascal Users Group, 2903 Huntington Rd., Cleveland, Ohio 44120. The Pascal Newsletter is a direct benefit of membership in PUG.

Membership dues in PUG are $\$ 25.00$ US regular, other forms of membership please inquire. Inquiries regarding membership should be sent to the above address. Newsletter correspondence and advertising should be sent to the editor at the aforementioned address.

Advertising Rates: $\$ 300.00$ Full Page. Please give your preference of magazine location: front, center, or back.

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Hello,
Well, this is the third issue I am involved with and there have been many changes. I would like to write of Pascal first.

Pascal has enjoyed a jump in attention in the last year. One reason is that there are Pascal compilers available for many machines and, I am tempted to say, they are available for any machine. Most of the major main frames have Pascal either directly or from a third party.

One step down in size, I know of only one machine, the Tandem computer which is without a Pascal implementation. A Tandem representative here in Cleveland informed me they have a language called "TAL" and in many cases will execute a Pascal program with no changes.

A couple more steps down in size are the small Digital Equipment machines and compilers are available from about four sources. IBM has the Display writer and Datamaster. These were released without our language, but in the last year, UCSD Pascal has been made available through IBM. Apple Computer has been a strong and long supporter of Pascal. TRS 80 has UCSD Pascal.

The smallest machine with Pascal is the TI 99/4A. In this size, Commodore has promised Pascal for this summer on the " 64 " and " 128 " machines.

The small computer, that is, the home computers and small business computers, have exceeded $\$ 10$ million in sales. This is according to Future Computing, a Richardson, Texas research firm.

With a guess, I would say that Pascal is implemented on at least $25 \%$ of these machines. If only $1 \%$ of these were being used to learn and program Pascal, then 25,000 people are presently involved. This is a lot of people looking for the best books from which to learn.

I am making an appeal to our members to submit comments and reviews of text books so that we all may benefit from your experience. I get calls from authors requesting information on Pascal. To these people, the best I can do is to send complete sets of Pascal News! With your comments and criticism, perhaps we could influence future text books.

Herb Rubenstein of Budget Computer in Golden, Colorado has sent a small article from Popular Computing. It seems that advanced placement test in computer science will use structure programming and the Pascal language. These tests allow up to one year of college level credits in computer science. The author of this article, Dan Watt, believes that the choice of Pascal in the testing may lead to Pascal as a defacto standard in high schools preparing students for college. Let me quote the last paragraph:
"This situation illustrates the power of the testing establishment to influence the lives of students and teachers. Although the vast majority of high schools now offer Basic as the standard computer language for most programming and computer science classes, this action by the College

Board may lead to the establishment of Pascal as a defacto standard for high school teaching and spawn an entire mini industry of curriculum to meet the new requirements. It may also offer significant school marketing advantages to microcomputer companies that already support Pascal - such as Apple, IBM and Texas Instruments."

I would like to see comments from you regarding this use of Pascal in a rite of passage.

In this issue, you will find a reprint of Dr. Srully Blotnick's column from Forbes magazine. I like this column because of the clever way he has made our economy dependent on you learning Pascal.

I enjoy Forbes magazine. They emphasize common sense and illustrate proven business practices. Forbes also takes a pulse of industries, and small computers is a fast growing industry. In a column called "'Technology", edited by Stephen Kindel, on March 28,1983 , he noted that $2 \%$ of the households in the U.S.A. own computers of one form or another. There had been predicitions of $40 \%$ of households by 1990 . This has been reduced to $20 \%$ in 1990 because there doesn't seem to be software that is useful in households.

Mr. Kindel ends this article with a quote from Seymour Papert, an MIT professor:
"The real purpose of learning how computers work should be to improve human logic and thought processes, to make people more creative, not simply more dependent on machines."

Maybe this would be a good issue to review the tools available in our back issues. This issue contains the APL scanner. I am embarrassed to print this, not because of the program's quality, but because it was submitted four years ago. Well, no time like the present.

In issue \# 17 (yellow), Arthur Sale submitted "Referencer'", a procedural cross reference. This program provides a printout of the heading of each procedure and function with indentation showing nesting. In issue \#25, Mr. Yavner has improved on this program with "A Better Referencer". Mr. Yavner claims that Pascal News has been his sole source of instruction in Pascal. I believe this is a compliment to Andy Mickel and Rick Shaw for their efforts to maintain this newsletter. We should also thank our contributors, Mr. Sale for instance, for outstanding generosity. These people will appreciate your complements, criticisms and gifts of money. (Ho! Ho!)

Andrew Tandenbaum, in issues 21 and 22/23, provided us with "The EM1 Compiler". This is a good look at all that is necessary for a pseudo 32-bit machine pascal compiler.

The UCSD Pascal Project started with a 16 -bit pseudo machine portable compiler. It was called P4 out of Zurich, Switzerland by Vrs Ammann, Kesav Nori and Christian Jacobi. I mentioned this because it has been published with critical commentary by S. Pemberton and M.C. Daniels in 1982. It is presented as a
case study of compiler design and is very interesting to read.

## Pascal Implementation

S. Pemberton and M.C. Daniels

Ellis Horwood Limited Publishers
Distributed by:
John Wiley \& Sons
605 Third Avenue
NY, NY 10016
USA
In \#21 you will find Jeff Pepper's fine implementation of extended precision arithmetic.

Nicklaus Wirth, Pascal's creator, wrote Pascal S and we have it in \# 19 (mislabeled \# 17). This is a subset of Pascal and was intended as a teaching aid.

Also in \# 19 is a Lisp interpreter written in Pascal.
"MAP", a Pascal macro preprocessor for large program development, is published in \# 17.

Issue \# 16 contains the Validation Suite version 2.2. This is the compiler checker that Arthur Sale and Brian Wickman have now revised to version 3. This new version is available by using the Validation Suite coupon in the rear of this issue.
"'Prose", a text formatter, by John Strait is the major program available in \#15. A disclaimer in the instructions manual admits that it doesn't do everything, but I must say, it has a lot of capability.

In \# 13, two programs were printed that performed the same work. A sort of "Battle of Algorithms". "Pretty Print" and 'Format'" used any Pascal programs as input and printed it in a consistent style.

For those of you looking for other Pascal periodicals, there are four of which I know. 'Pascal Market News'', 30 Mowry Street, Mt. Carmel, CT. 06518. This is a nice quarterly for $\$ 9$.

Another quarterly for Oregon Software users is the '"Pascal Newsletter". Maybe this is too narrow in content, but you will know what Oregon Software is up to. Their address is 2340 SW Canyon Rd., Portland, OR. 97201.

A very slick magazine with good design is "Journal of Pascal and Ada." You can contact them at West Publishing Company, 898 South State Street, Orem, UT. 84057. The cost is $\$ 14$ for six issues.

The USUS News and Report is more a system user's journal, but the system is based on Pascal. They also have a software library, seventeen floppy disks full, and all in source code and written in Pascal.

Now to the business of Pascal News. Pascal News, as the Pascal periodical granddaddy published since January 1974, has had its ups and downs. In 1979 our circulation was 7,000 ; now it is 3,600 . Our biggest problem has been irregular publication. I am committed to four issues this year and I am considering six issues next year. I believe that regularity will supply us with growth and members and more software tools.

As I mentioned in the last issue, PUG (AUS) has stopped and I, in the USA, have taken over their area. Unfortunately, they have not sent me their mailing list and I fear that I have lost touch with our members there. This issue will be sent to those members listed as of 1979 and I hope they will 'spread the word' and the subscription coupons!

Our PUG (EUR) has performed very nicely and I thank Helmut Weber and friends for their good work. But they have a problem concerning money. They have not charged enough for subscriptions and were pressed to send our \#24. As a result, I will mail all issues directly and I hope you will not be inconvenienced. Please keep in touch with them as they are a strong group.

I have saved the worst for last. In November, 1982, I sent 300 copies of issue \#24 to Nick Hughes in care of PUG (UK), Post Office Box 52, Pinnen, Middlesex HA5 3FE, United Kingdom. Using the phone number 866-3816, the air express shipper delivered these issues by mid-November. All well and good. The issues arrived before the cover date with plenty of time to post them to our English members. I called Nick at this number many times, but spoke to him only after many months. It was late April and I asked if I should use the same procedure in shipping \#25 to him.

Nick said that the issues arrived properly and that method was efficient but wanted to know what was in \#25. He told me that he did not like issue \#24 and from the sound of it, did not like issue \#25. He had disliked \#24 so much, he decided not to send any of them out. Need I say more?

Nick will not supply his mailing list so I am sending this issue and \#25 directly to the members of record in the United Kingdom as of 1979. If you feel a need to find out why Nick Hughes did not like issue \#24 or would like to see it yourself, please call or write Nick at the above address and ask for your copy. He has 300 and I am sure he can spare one.

As a result of these difficulties, I will receive and service all subscriptions from here in Cleveland, Ohio. From now on, there will be only one person to blame if you have a complaint.

As of this issue, a year's subscription is raised in price to $\$ 25$ a year and $\$ 50$ for three years. These represent two sets of costs; production and organization. Production costs are typesetting, printing and mailing. Other activities of production are editing, reviewing, quality assurance and formatting. These tasks are performed by "yours truly" and presently I do them for free. (I'm real smart!)

Organization is a cost of servicing you and other members satisfactorily. This includes collecting and reviewing the mail, depositing checks, updating the mailing list, sending back issues to fill new subscriptions and sending sets of previous years back issues. In order to do this correctly, and in a timely fashion, I don't do it. I pay a firm to perform 'fulfillment'" and it takes one or two days per week. This cost is small compared to the bad feelings generated if not done correctly and quickly.

These are costs of which you are totally responsible. This newsletter has been a beneficiary of volunteerism. There are no volunteers now (save me). In many magazines, advertisements will pay for all production and organizational costs plus provide profits, sometimes large profits.

The cost of a full page ad in Byte or PC or PC World is over $\$ 2,000$ and these are publications with 500 pages !

Now we may be able to keep our costs down and publish more often if we accept advertising. Three hundred dollars per page is not expensive. I will pursue
advertisers and I am asking for your help. If you are writing a book, have your publisher advertise with Pascal News. If you are making software packages, influence your boss in the virtues of an ad in Pascal News. If you manufacture or sell computers, sell your product from the pages of Pascal News. This is the oldest Pascal publication and, I proudly say, the most influential.

This newsletter help spread Pascal and our members were most influential in the standard efforts.

I believe Pascal News' new mission is to enable Pascal to be taught in the easiest way. This is in many forms. For instance, reviews of books and texts, discussion of what features to teach first as a foundation, how to teach advanced courses, discussions of extensions or standard program tools to include in every well written program as it is appropriate.

By the way, Andy Mickel tells me that the "Pascal User's Manual and Report'" by Jensen and Wirth has sold 150,000 copies in 1982 . This is interesting considering that in the previous seven years, it sold 175,000 copies. A very sharp jump in interest.

A new text book has been sent to me, "Pascal"' by Dale/Orshalik, 1983 DC Heath. A nice title, short and to the point. The preface states a philosophy that I would like you to comment on.
"In the past there have been two distinct approaches used in introductory computer science texts. One approach focused on problem solving and algorithm design in the abstract, leaving the learning of a particular language to a supplemental manual or a subsequent course. The second approach focused on the syntax of a particular programming language, and assumed that the problem-solving skills would be learned later through practice.

We believe that neither approach is adequate. Problem solving is a skill that can and should
be taught - but not in the abstract. Students must be exposed to the precision and detail required in actually implementing their algorithms in a real programming language. Because of its structured nature, Pascal provides an effective vehicle for combining these two approaches. This book teaches problem-solving heuristics, algorithm development using top-down design, and good programming style concurrently with the syntax and semantics of the Pascal language."
One of the letters mentions high resolution graphics. I know of two texts that use Pascal as the illustrative language of their algorithms. They are "Principles of Interactive Computer Graphics" by Williams Newman and Robert Sproull, 1979 McGraw-Hill and "Fundamentals of Interactive Computer Graphics" by James Foley and Andries Van Dam, 1982 Addison-Wesley.

Two notes from members:
Steven Hull of Campbell, California, received a notice from me that \#22/23 had been returned to us because the postal service will not forward bulk mail. His reply:
"I guess this will teach me to move from Lakewood (a suburb of Cleveland, Ohio). Didn't know bulk mail wasn't forwardable. The Postal Diservice has been re-routing every piece of junk mail for a full year . . . I might have to file suit to stop it all!

And from Eric Eldred of New Hampshire who rewarded Pascal News with a three year subscription and dutifully filled the coupon with name and address and arrived at a request for "Date". Eric filled in "No! Married!". Thanks Eric, I needed that!

To Charlie Gaffney,
I'm glad you have taken on Pascal News. I hope it works.

Perhaps, I should say what I would like to see published in Pascal News. The most valuable things are 1) Tools, and 2) Info on the various implementations.

In my job we are using many computers. It is very helpful to know which compilers work well, meet standards, and produce efficient code. Apple Pascal is nearly bug free, and works as specified (with UCSD quirks). IBM Pascal VS is good - extensions are large presenting conversion problems if they are used. It has a good interface to FORTRAN. VAX Pascal is plain vanilla, appears to work well but we have not tested it in difficult situations. HP Pascal 1000 works fine but does not have a stack architecture and seems to compile slowly. Recent tests on HP Pascal 1.0 for the HP 200 computers seem to indicate it derives from UCSD although it is a native code 68000 compiler. It seems to work very well. We are interested in Pascal for the Data General Eclipse.

> Good luck, Dennis Ehn
> 215 Cypress Street
> Newton Centre, MA 02159

## Gentlemen:

Would you be so kind as to send information on the Pascal User's Group (PUG) and its official publication Pascal News. Recently we have acquired a microcomputer Pascal compiler and are very much interested in keeping up with current developments in Pascal.

Our system is based upon a SouthWest Technical Products Corporation S/09 computer, running the UniFLEX Operating System (similar to UNIX). If specific information is available for this unit, please let us know.

Additionally, the college has several (approximately 18) Apple computers which are capable of running the UCSD Pascal System. Once again, any special information here would be very helpful.

We look forward to hearing from you and hope that we can make a positive contribution to the Pascal User's Group.

> Yours Truly Lawrence F. Strickland Dept. of Engineering Technology St. Petersburg Jr. College
> P.O. Box 13489
> St. Petersburg, FL 33733

Pascal compiler. What makes it surprising is that to the best of my knowledge I have never sent in an entry, and the information provided is about a year and a half out of date.

In case you would like to provide your readers with valid information, I have enclosed an implementation note for the currently available compiler. I have also enclosed a copy of the ISO validation suite report from our language manual.

Work is currently being done on moving this compiler to the 68000 family of processors and should be available by the end of 1983.

On another note, I have received issues number 21, $22 / 23$, and 25 , but not issue 24 . I am also enclosing a check for a 3 year membership - please see if you can determine what happened to number 24.

Sincerely, Robert Reimiller Owner, OmegaSoft 5787 Brandywine Ct. Camarillo, CA 93010

December 1, 1982
I hope the letter referring to the possible end of the P.U.G. is wrong! I can be of some help if needed.

> Allen Duberstein Pine Instrument Co. 3345 Industrial Blvd. Bethel Park, PA 15102

January 10, 1983
Dear Mr. Gaffney:
Enclosed is a check covering both the remailing cost of Pascal News \#24 (\$5) plus my membership renewal for two years (\$18).

My apologies for getting out of synchronization with the Pascal Users Group. As the post office informed you, I recently moved to the address noted. Frankly, I hadn't received a Pascal News in so long that I simply forgot about it. It appears that I won't miss any issues - the enclosed All-Purpose Coupon is from issue \#23.

Interestingly, after a long period (3 years) of not using Pascal, it looks like I will be using it once again. We have a couple of Convergent Technologies workstations in my office. These are very nice 8086-based machines; Burroughs sells them as the B-20s, and NCR sells them as WorkSavers. We will probably be getting a Pascal compiler, and I am looking forward to getting back into Pascaling in the near future.

November 30, 1982
I was surprised and pleased to receive issue number 24 of Pascal News. Thanks for taking it over. I do have one question, however, which you might be able to help me with. What year is it? My address label includes [82] on it but the previous issue I received was dated September, 1981. I notice that this issue is dated January, 1983. Should I send in another year's subscription money now? What happened to 1982 ? I never have managed to figure our Pascal News' subscription scheme. Maybe a note in the issues towards the end of a year saying "if your address label says [82] it's time to send in a renewal'' would help.

Thanks for your help.

Richard Furuta<br>Computer Science, FR-35<br>University of Washington<br>Seattle, WA 98195

## 8 February 1983

Dear Sir,
I received your notification of renewal in the mail yesterday. I am slightly concerned that you may not have received the check which I mailed to you in December. I hope that it has only been a slight mix-up, and in fact, my subscription has been renewed for 3 years, as I requested.

I am currently using the Pascal implemented by Microsoft for the IBM Personal Computer. It has some non-standard features which were provided in order to allow programmers to access the full capabilities of the machine. This implementation is quite flexible, and was designed to allow users to produce systems programs, as well as application programs.

The greatest shortcoming to this product, however, is its lack of usable documentation. Even someone like myself, who has been programming in Pascal for 8 years, has difficulty in trying to locate the appropriate material in the 'reference manual'. Once this is overcome, the user is able to use this version for the production of some very powerful software.

I continue to look forward to the delivery of your fine newsletter. I enjoy the articles, and realize how difficult a task you have. Keep up the good work.

Regards,
Robert A. Gibson 1609 Lake Park Dr. Raleigh, NC 27612

November 30, 1982
Pascal is being used for process control of laser trimming systems. We use Oregon Software Pascal.

Barbara Huseby, Training Dept.
Electro Scientific Industries 13900 N.W. Science Park Drive Portland, OR 97229

March 3, 1983
Dear Mr. Gaffney:
I'm writing to let you know why I am not renewing my subscription to Pascal News. The main reason is that the price is now too high for the utility of the product (at least to me). I appreciate your efforts to keep PUG and Pascal News going, but I'm afraid they may have outlived their usefulness. Pascal is not really in need of promotion as it was when PUG was formed. The Journal of Pascal \& Ada may be an appropriate successor.

As a long-time subscriber and occasional contributor, I wish you luck in your efforts.

Richard Leklanc<br>Assistant Professor<br>Georgia Institute of Technology<br>Atlanta, GA 30332

January 7, 1983
Hang in there, Charlie!
Andy Mickel
106 SE Arthur Avenue
Minneapolis, MN 55414

## December 9, 1982

Dear Sirs:
Could you provide us with information on membership in your organization, both personal and institutional, as well as the subscription cost of your journal.

We are also interested in a rigorous comparison of the various PASCAL versions implemented by mini and microcomputer vendors. Do you know of any such comparative research? We are making plans to offer Advanced Placement Computer Science in the fall term of 1983, and wish to select an effective computer.

Very truly yours,
Charles McCambridge
Director
Instructional Materials Services
Niskayuna High School
1626 Balltown Rd.
Schenectady, NY 12309

December 25, 1982
Merry Xmas! Good luck, Charlie! Is your "acquisition", of PUG a sign that PUG and USUS will someday merge? I'm not sure I'd like that, but let's see.

Jim Merritt
P.O. Box 1087

Morro Bay, CA 93442

December 24, 1982
Please send me information on joining the Pascal

User's Group, I am a software project engineer at General Electric in Syracuse. I am currently in the process of selecting a high level language for internal programming of a $1024 \times 1280$ resolution raster display. Pascal is the leading candidate, therefore, I am very interested in the latest information regarding the language which I feel a user's group could provide.

My interest does transcend my work however as I do own a Commodore SuperPET which includes the University of Waterloo software package consisting of Pascal, APL, Fortran, Basic and a 6809 Assembler.

Sincerely,<br>Douglas W. MacDonald 4303 Luna Course Liverpool, NY 13088

## 2/5/83

## To Whom It May Concern:

I just received your notice to inform me that my membership is about to expire and that I should renew now.

I would like to tell you that I would consider renewing if I could be assured of getting my money's worth - this time!

When I first joined in 1981, I didn't hear from Pascal News for almost a year. Then a few months ago, I received a second issue, but that's been it.

Now I am a convicted Pascaler. I understand the difficulties of operating a non centralized club, but \$20 should buy some kind of organization for things I feel.

Can you assure me of a better value this time around?

> Cordially,
> David Abate
> Micro People 116 S. Bowdion St. Lawrence, MA 01843
P.S. Question: Do you intend anything on UCSD-Pascal? This is my greatest interest.

7 January 1983
Hi,
This is a note in a bottle to: 1 ) find out if you're still out there, and 2) what's happening with Pascal. It doesn't seem to be taking the bite (or is that byte) out of Basic I thought it would.

We will start covering Pascal as soon as we have finished Basic programming - about five weeks from now. The extension program from Hocking Technical College in Nelsonville has provided seven Apple II and Apple III computers and two printers. By the end of the year, they will have installed a winchester disc and either a modem or a microwave link to their main campus computer. We'll need it by then to cover the Cobol and Fortran IV programs we'll be writing.

Most of my practical computer experience is in assembler language. I used it at Cincinnati Milacron's

Process Controls Division (Mater's of the controls for the $\mathrm{T}^{3}$ Industrial Robot).

I am interested in any literature you have to send me. In particular, I would like the titles of the books you consider best for teaching Pascal - either on the Apple II or on computers in general. Apple, Inc., sent me the Pascal Reference Manual (just a bit or a nibble over my head). I've also read copies of the DOS 3.2 Reference Manual and their Basic Programming Manual. I covered all these before classes started and wound up tutoring two other student/inmates.

Sincerely,
Brian Appleman 166-767
15802 St. Rt. 104
P.O. Box 5500

Chillicothe, OH 45601
P.S. If you need more on my background, just ask.

## 83-02-24

Dear Charlie:
I am a member of PUG (AUS) which has just folded, and I would like to re-enroll through PUG (US).

I don't share Arthur Sales view that PUG and PN have no purpose now that there is an ISO standard. The world still needs cheap, good software and PN (in a modest way) supplies some of it. Also, some organization is needed to defend and develop good programming language and style.

PUG (AUS) says I have a credit of 12 (old) issues and that the funds have been sent to you. Please will you accept my re-enrollment and advise me how many (new) issues I am now entitled to?

Finally, I, and I'm sure, many others appreciate your offer to keep PUG/PN going.

Thanks again.
Yours sincerely,
Peter Edwards
40 Davison St.
Mitcham, Vic.
Australia 3132

## December 3, 1982

Best wishes in this venture, Charlie. I agree that Pascal News and P.U.G. are worth saving.

John W. Baxter
750 State Street, Apt. \#224
San Diego, California 92101

## February, 1983

You people have ripped me off for the last time!
By your own back order form (attached) you show that my renewal in 1981 paid for 3 issues mailed in 1982. But then, WHAT OF MY RENEWAL PAID IN 1982? ONLY ONE ISSUE \# 24 COUNTS??? AND THAT HAD TWO PREVIOUSLY PUBLISHED PRO-

GRAMS!! (That is, programs I had ALREADY received.) If you ran a decent organization, you'd make my 1982 renewal count for 1983 also.

David S. Bakin<br>Softech Inc.<br>360 Totten Pond Road<br>Waltham, MA 02154

December 24, 1982
We're indebted to you, Charlie!

> Wayne N. Overman 3522 Rockdale Ct. Baltimore, MD 21207

February 17, 1983
Dear Mr. Gaffney,
I am one of those folks who does not have a currently correct address with Pascal News.

Enclosed is a check for $\$ 5$ for a copy of issue 19 which was returned to you.

Thank you on behalf of all the members of the user's group for the effort you are putting out. It is very much appreciated.

Tom Bishop
P.O. Box A

Kenmore, WA 98028

March 14, 1983
Dear Sir or Ms.:
We plan to offer Pascal at our school. I would appreciate receiving information on your group and, if possible, a sample copy of Pascal News.

Any suggestions or information you could send would be appreciated. We are particularly concerned that the new Apple 2-E does not support Pascal with one disk drive. We had hoped tht UCSD Pascal with one drive would work on the Apple 2-E.

Thanks for your help.
Sincerely,
Harold Baker
Director, Computer Science
Litchfield High School
Litchfield, CT 06759
February 11, 1983
Hi !
Here's my renewal. I really enjoy Pascal News and have been upset about what has happened with it the past 18 months or so. It has been of substantive value to me, particularly in the area of the style of Pascal coding among the community that have submitted articles.

I would like to see more articles on Modula 2,

Wirth's follow on to Pascal and Ada in parallel. To me, this would seem a way of keeping PUG alive as well as providing a growth path to these languages for Pascal programmers.

I use Pascal/VS extensively at work and I have found its extensions the best of any other Pascal compiler for $S / 370$ compatible machines. Almost all of its extensions are within the "spirit" of Pascal and uses a very good extension to STRING data. Of particular convenience is its READSTR and WRITESTR functions (they are procedures actually -unfortunately). I force the concept of function upon them by embedding their invocation within a function when required.

I never received issues 20 and 21 of Pascal News during the confusion, although I did mention this at times. I would certainly purchase them separately, but I am not prepared to purchase two sets to get them. Please advise.

Thanks for your work,
Bob Dinah
630 Alvarado St. \#207
San Francisco, CA 94114

## November 12, 1982

Dear Pascal User's Group:
The only source of information that $I$ have on the Pascal User's Group came from "The BYTE Book of Pascal'", according to an article written by Kenneth Bowles. An editor's note of July 1, 1979 listed the annual newsletter as $\$ 6.00$ per year. I am enclosing $\$ 12.00$ in case things have increased since that date. If this amount is insufficient, please make it up on back issues.

I am currently using an Apple /// with Apple computer's version of UCSD Pascal. There does not seem to be more than a dozen books written on Pascal, and just a few on UCSD.

I am an ex-electrical engineer, turned to building construction. Previously, I worked for Westinghouse Research Center in Pittsburgh, and used the Burroughs B6500 main frame computer with ALGOL language. The B6500 used a number of formats and types that I miss; the Fixed Format was especially useful since it allowed the user to specify the number of total digits and the number of decimal digits combined. I would like to use this format in UCSD Pascal.

Thanks for taking the time to help me.
Very truly yours,
Larry J. Moorhead
5207-32nd Street East
Bradenton, Florida 33508
18 March 1983
Dear Sirs,
For the first time we have received a copy of Pascal News, and it has been read with great interest.

We would like to join your User Group but cannot find either a price or contact address for our region.

Please send us this information as soon as possible,
so that we can become members and start receiving your journal on a regular basis.

We have taken note of your abhorrance of paperwork (and endorse the sentiment) and will send the necessary prepayment once we receive the information.

Yours sincerely,
Bette Kun
Librarian
Control Data
P.O. Box 78105

Sandton, South Africa 2146

20 April 1983
Dear Mr. Shaw:
Enclosed is a check for $\$ 10.00$ for a one-year subscription to the PASCAL Users' Group Newsletter. We have just recently acquired PASCAL-2 here at Villanova and our students are using it on LSI-11 systems running RT-11 V4.0 for applications involving real-time control, data acquisition, and computer communications.

Sincerely yours,
Richard J. Perry, Ph.D. Villanova University
Dept. of Electrical Engineering Villanova, PA 19085

15th February, 1983
Dear Mr Gaffney,
As a long PUG user the demise of PUG-AUS is a blow. Anyhow, as you can see from the attached letter I would love to continue and thus need your help.

Could you please detail the fees for 1983 for us "down under" for surface mail and air mail and as you can see I'm afraid I've not got issue number 21. Can you help?

For interest I use:

| UCSD Pascal/p-System | ERA-50 Computer <br> (8-bit, 8085 base) |
| ---: | :---: |
| Pascal MT + under CP/M | ERA-50 Computer |
| (8-bit, 8085 base) |  |
| and MP/M | ERA-50 Computer |
|  | (8-bit, 8085 base) |
| al MT+86 under CP/M-86 | ERA-80 Computer |
|  | (16 bit, 8086/8087 based) |
| and MP/M-86 | ERA-80 Computer |
|  | $(16$ bit $8086 / 8087$ based) |

Regards,
Dr. William J. Caelli, F.A.C.S.
President
ERACOM Group of Companies
P.O. Box 5488, G.C.M.C. Qld. 4217, Australia
me introduce myself first. I am a student pursuing a computer course in the Hong Kong Polytechnic - a licensed user of your OMSI-PASCAL-2 V1.2. I don't know what your definition of user may be. May it be my Polytechnic or any student or programmer who use your OMSI PASCAL-1 under the Polytechnic, I venture to call myself a user in this letter, and would like to join the Pascal Users' Group and receive the newsletter.

In the past few months, I have been doing extensive programming using PASCAL, and find it very handy, especially in writing structured programs. However, until recently when I develop some system programs, I find problems. I discover that there is no source listing or documentation on the OMSI PASCAL-1 run time system (possibly in file FPP.RTS) and its relationship with RSTS/E, and I cannot interface with the low level I/O trap handlers without knowing their details. I find some problems on the RESET ODT mode, but I cannot deal with it in assembly level.

All in all, my problem is highly personal and does not in any way bear relation with the Hong Kong Polytechnic. However, as a student on computing, I don't want to leave problem unsolved. So, please send me any informational help, if possible.

Included please find a bank draft of $\$ 6$ for subscription.

May I state once more my request. I need information on OMSI PASCAL-1 run time system especially the EMT trap handling.

Thank you very much in advance.
Yours faithfully,
Mr Kam Man-Kai
Flat 8
3/F Ting Yin House
Siu On Court
Tuen Mun-N.T.
Hong Kong

6th November, 1982.
Dear Mr. Mickel,
I am a student of computing studies in the H.K. Polytechnic. Recently, I got a chance to buy a Chinese version of 'A Practical Introduction to Pascal' by Wilson \& Addyman from which I was informed that there is a PUG in States.

Briefly understanding the objectives of the PUG, I find myself in great interest in joining the group. Would youb be so kind as to provide me with further information as far as the PUG is concerned. I am eagerly looking forward to your reply.

Yours sincerely,
Alan Kwong
12, Boundary St. Po Hing Bldg.
8/F, Block 'C'
Kln., H.K.

Dear Sir/Madam,
This is the first letter I write to contact you. Let

## December 23, 1982

We have been using Oregon Software's RT-11 Pascal implementations for over three years with excellent results and complete satisfaction; Pascal is used for scientific 'number crunching', program development, algorithm testing, etc.

Bob Schor The Rockefeller University 1230 York Avenue New York, NY 10021

December 30, 1982
A worthwhile journal.

> George Williams
> Union College Schenectady, NY 12308

## February 3, 1983

Good job, Charlie! and good luck to the renewed Pascal News!

Norman W. Molhant 320 Principale<br>Tres-Saint-Redempteur, P.Q.<br>Canada J0P 1P0

May 1, 1983
A professor in Ithaca, NY told me there exists a public domain UCSD Pascal available for micro's.

I have a 60 K Z- 80 which uses memory map video, and a $63 \mathrm{~K} 8085 / 8088$ (both machines $\mathrm{S}-100$ bus) which uses a TVI 950. I also have a H-29 terminal (like Z-19 but with a detached keyboard).

Is there really any way of getting this UCSD Pascal running on one of my systems? (I have UCSD on the Sage also Modula-2. Good stuff.)

Thanks,
J. E. Pournelle, Ph.D.

12051 Laurel Terrace
Studio City, CA 91604

Yours sincerely,
M.J. Palmer

CSIRO
Private Bag
P.O. Wembley, W.A. 6014

## Program APLscanner

## By Vincent Dichristofano, Alan Kaniss, Thomas Robinson, and John Santini NADC, Philadelphia, PA

```
program APLscanner(input ( + TERMINAL }, output . APLfile );
84
| Purpose:
    This program is an impienentation of APL in Pascal.
* Authors:
        Vincent Dichristofano
        Alan Kaniss
        Thomas Robinson
        John Santini
            authors' affiliation - NADC
                                    Phil. PA. USA
        project leader: Dr. Joseph Mezzaruba
        This program was written as part of an independent study
        course at Villanova University.
* Submitted and accepted for Pascal News. DEC 197%.
label
\abel
const
    prefix1 = 60;
    prefix2 = 62'{ prefix for CDC ASCII 12-bit codes };
    MaxVarNameLength = 10;
    MaxINputLine = 132;
    InputArraySize = 134;
    NumberOfMessages = 100;
    MessageLength = 80;
type
    PackedString = packed array [1 .. MaxVarNameLength] of 0 .. 8191;
    TokenNoum =
        (FormRes, FormArg, GlobVar, MonadOper, ReductOper, DyadOper,
        SpecOper, constant, StatEnd);
    values = record
                    Realval: real;
                    Nextvalue: `values
                end;
```



```
            VarName: PackedString { v1 };
            FuncTabPtr: `FuncTab { v2 - ftab };
            ValTabPtr: `ValTab { v3 - vtab };
            DeferedValTabPtr: "FParmTab;
            NextVarTabPtr: "VarTab
                end;
    ValTab = record
            \mathrm{ IntermedResult: Boolean;}
            dimensions: integer;
            FirstDimen: 'DimenInfo;
            ForwardOrder: Boolean;
            FirstValue: "values;
            NextValTabLink: `ValTab
            end;
    TokenTable record
                    NextToken: "TokenTable;
                    case noun: TokenNoum of { p } }
                    (VarTabPtr: `VarTab);
                MonadOper: (MonIndex: integer);
                ReductOper: (RedIndx: integer);
                DyadOper: (DOpIndx: integer);
                SpecOper: (CharIndx: integer);
                constant: (ValTabPtr: `ValTab);
                StatEnd: (EndAdj: integer)
            end;
    vfunc = record
            NextStmnt: "TokenTable;
            NextVFunPrt: "vfunc;
            StatLabel: PackedString
            end;
    OperatorTypr = (niladic, monadic, dyadic);
    FuncTab = record
            FuncName: PackedString { f1 };
            FuncName: PackedString { f1 }
            arity: Operatoryype f f% frue = explicit };
            result: Boolean { f3 true = expl
            ResultName: PackedString { f4
```



| 371 | TokenError := true; | 473 | $\underline{v a r}$ |
| :---: | :---: | :---: | :---: |
| 372 | for MsgCol := 1 to MessageLength do | 474 | index: integer; |
| 373 | Write(ErrorMsgs[ErrorIndex, MsgCol]); | 475 |  |
| 374 | writeln; PrintAPLStatement ( echo statement to user ); | 476 | begin 1 see if two names (identifiers) are the same |
| 375 | for MsgCol $:=1$ to(position-1) do write(' '); | 477 | NamesMatch := true; |
| 376 | Writeln(chr(character[UpArrow])) ( print pointer to user error ); | 478 | for index := 1 to MaxVarNameLength do |
| 377 | end \{ error \}; | 479 | if NameOne[index] <> NameToo[indē] then NamesMatch := false |
| 378 |  | 480 | end $T$ namesmatch $\}$; |
| 379 |  | 481 |  |
| 380 | procedure SkipSpaces; | 482 |  |
| 381 |  | 483 | procedure TableLookUp(Testchar, Tablelength: integer; table: OpTable; |
| 382 | beg in | 484 | var TableIndex: integer); |
| 383 | while (APLstatement[position] = character[space]) and (position <= | 485 |  |
| 384 | LineLength) do | 486 | var |
| 385 | position := position + 1 | 487 | Index: integer; |
| 386 | end \{ skipspaces \}; | 488 |  |
| 387 |  | 489 | begin $\{$ check for membership in a given table |
| 388 |  | 490 | TableIndex : = 0; |
| 389 | procedure GetAPLstatement; | 491 | for index := 1 to Tablelength do |
| 390 |  | 492 | if TestChar $=$ table[index]. OpSymbol then TableIndex $:=$ index |
| 391 | var ${ }^{\text {In }}$, | 493 | end $T$ tablelookup \}; |
| 392 | Input Char: char; | 494 |  |
| 393 | TestforPrefix: integer; | 495 |  |
| 394 | FirstTry: Boolean; | 496 | procedure identifier(var name: PackedString; var ItsAnIdentifier: |
| 395 |  | 497 | Boolean) ; |
| 396 | $\frac{\text { begin }}{\text { for LineLength }}:=1$ to MaxINputLine do | 498 |  |
| 397 | APLstatement[LineLength] $:=$ character[space] \| blank out line \}; | 499 | $\frac{\mathrm{Var}}{\text { NameLength: integer }}$ |
| 399 | LineLength $:=0$; FirstTry $:=$ true; position $:=1$; | $\begin{aligned} & 500 \\ & 501 \end{aligned}$ | NameTooLong: Boolean; |
| 400 | LineTooLong $:=$ false; ${ }_{\text {APLstatement }}$ [InputArraySize] $:=$ character[omega]. | 502 |  |
| 401 | APLstatement[InputArraySize] : $=$ character[omega]; | 503 | beg in |
| 402 | APLstatement[InputArraySize - 1] : $=$ character[space] | 504 | ItsAnIdentifier := fatse; SkipSpaces; |
| 403 404 | $\underset{\text { repeat }}{ }$ set end-of-line ${ }^{\text {a }}$ | 505 | if ItsALetter(APLstatement[position]) |
| 404 | $\frac{\text { repeat }}{\text { begin }}$ | 506 | then |
| 405 406 |  | 507 | NameTooLong := false; ItsAnIdentifier := true; |
| 406 | FirstTry := false; | 508 | NameTooLong := false; ItsAnIdentifier := true; <br> for NameLength := 1 to MaxVarNameLength do \{ blank out nam |
| 407 | while (not eoln(input)) and (not LineTooLong) do | 509 510 | name[NameLength] : = Character[space] - |
| 409 | if Linelength < MaxinoutLine | 511 | NameLeng th : $=0$; |
| 410 | then | 512 | while (ItsALetter(APLstatement[position])) or (ItsADigit ( |
| 411 | begin | 513 | APLstatement[position])) do |
| 412 | LineLength := LineLength + 1; read(InputChar); | 514 | begin \{ build identifier \} |
| 413 | The following code would be removed for non-CDC installations | 515 | NameLength := NameLength +1 ; <br> if Namelength <= MaxVarNameLength |
| 414 | TestforPrefix := ord(InputChar); | 516 | then name[NameLength] $:=$ APLstatement[position] |
| 415 416 | $\underline{\text { if }}$ (TestforPrefix $=$ prefix1) or (TestforPrefix $=$ prefix2) | 517 518 | else NameTooLong := true; |
| 417 | then | 518 519 | position := position + 1 |
| 418 | begin | 520 | end; |
| 419 | read(InputChar); | 521 | if NameTooLong |
| 420 | APLstatement[LineLength] := 100 * TestForPrefix + ord( | 522 | then SError (70) i name greater than maxlength |
| 421 | end InputChar); | 523 | end ${ }^{\text {I }}$ identifier ); |
| 422 | elsend | 524 | end identifier ; |
| 423 | else | 525 |  |
| 424 425 | \{ \} | 526 527 | procedure MakeNumber(var RealNumber: real; var ItsANumber: Boolean); |
| 426 |  | 528 |  |
| 427 | end APLstatement[LineLength] := ord(InputChar) | 529 | sign, Digit Count: integer; |
| 428 | el $\frac{\text { end }}{\text { se }}$ LineTooLong : $=$ true | 530 | sign, Digitcount: integer; |
| 429 | end | 531 | begin \{ convert character input string to numerical representation |
| 430 | until Linelength <> 0 \{ reject null lines \}; | 532 | ItsANumber $:=$ false; SkipSpaces; sign $:=1$ 1; DigitCount $:=0$; |
| 431 | if LineTooLong then SError (71) | 533 | RealNumber :=0.0; |
| 432 433 | end ${ }^{\text {d }}$ getaplstatement ; | $534$ | if (APLstatement[position] = character[negative]) or (ItsADigit ( |
| 433 434 |  | 535 536 | -APLstatement[position])) |
| 435 |  | 537 | then ${ }^{\text {a }}$ |
| 436 | function ItsADigit(Test Char: integer) : Boolean; | 538 | $\frac{\text { beg in }}{\text { Its }}$ ANumer $:=$ true. |
| 437. | var | 539 | if APLstatement[position] = character[negative] |
| $438{ }^{4}$ | DigitIndex: APLCharSet; | 540 | then begin sign $:=-1 ;$ position $:=$ position + 1 end; |
| 441 | begin \{ test to see if input character is a digit | 541 | if not ItsADigit(APLstatement[position]) |
| 442 | ItsADigit $:=$ false; | 542 | then |
| 443 | for DigitIndex : = OneSymbol to Zerosymbol do | 543 544 | begin ${ }^{\text {SEr }}$ ( |
| 444 | if Testchar = character[DigitIndex] then ItsADigit := true | 545 | SError(1) \{ digit must follow a minus sign \}; |
| 445 | end $T$ itsadigit 3 ; | 546 | ItsANumber := false; <br> end |
| 446 |  | 547 | else |
| 447 |  | 548 | begin \{ form whole number portion \} |
| 448 | function ItsALetter(TestChar: integer) : Boolean; | 549 |  |
| 449 |  | 551 | white ItsADigit(APLstatement[position]) do |
| 450 | var | 552. | begin |
| 451 | LetterIndex: APLcharSet; | 553 | Real Number : $=10.0$ * Real Number + Char ToNum(APLstatement |
| 452 |  | 554 | [position]); |
| 453 | begin \{ test to see if input character is a letter \} | 555 | position := position + 1 |
| 454 | ItsALetter := false; | 556 | end; |
| 455 | for Letterindex : = asymbol to ZSymbol do | 557 | if ${ }_{\text {APLstatement[position] }}=$ character[period] |
| 456 | if Testchar = character[LetterIndex] then ItsALetter : = true | 558 | then |
| 457 | end \{itsaletter \}; | 559 | begin |
| 458 |  | 560 | position : = position + 1; |
| 459 |  | 561 | while ItsADigit(APLstatement[position]) do |
| 460 | function CharToNum(TestChar: integer): integer; | 562 | begin \{ form fractional portion \} |
| 461 |  | 563 | Real Number : = Real Number + Char ToNum(APLstatement[ |
| 462 | var | 564 | position) * $\exp ((-1.0-$ DigitCount) * 2.3025851 |
| 463 | Digit Index: APLcharset; | 565 | ); |
| 464 |  | 566 | DigitCount := DigitCount + 1; |
| 465 | begin chage a character to a number \} | 567 | position := position + 1; |
| 466 | for Digit Index : = OneSymbol to ZeroSymbol do | 568 | end; |
| 467 | if TestChar = character[Digit Index] | 569 | if Digitcount $=0$ then |
| 468 | then Char ToNum := digits[DigitIndex] | 570 | begin |
| 469 | end \{ chartonum \}; | 571 | SError(2) \{ digits must follow a decimal point \}; |
| 470 |  | 572 | ItsANumber : $=$ false; |
| 471. | function NamesMatch(NameOne, NameToo: PackedString): Boolean; | 573 574 | end; ${ }^{\text {end }}$ |



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\begin{array}{rl}{RealNumber := RealNumber * sign }\end{array}
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\begin{array}{l}{RealNumber := RealNumber * sign }\end{array}\mp@code{end}
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```
if (Newt
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if (Newt
```

procedure DyadicOpCheck;
$\frac{\text { var }}{\text { TableIndex: integer; }}$
$\frac{\text { var }}{\text { TableIndex: integer; }}$
begin
TableLookUp(APLstatement[position], 16, DOpTab, TableIndex);
if TableIndex $=0$
if Then
beg in
$\frac{\text { beg in }}{\text { Tabl }}$
TableLookUp(APLstatement[position], 12, CharTab, TableIndex);
if TableIndex $=0$
if Tab
$\frac{\text { if }}{\text { en }}$ APLstatement[position] $=$ character[SouthCap]
then

OldTokenPtr := SaveTokenPtr; dispose(NewTokenPtr);
$\begin{array}{lll} & & \\ \text { OLdTokenPtr }:=\text { SaveTokenPtr; } & \text { dispose(NewTokenPtr); } & 730 \\ \text { NewTokenPtr }:=\text { SaveTokenPtr; } & \text { position }:=\text { LineLength }+1 ; & 732\end{array}$
$\begin{array}{ll}\text { NewTokenPtr }:=\text { SaveTokenPtr; position }:=\text { LineLength }+1 ; 732 \\ \text { end t this was a comment - ignore remainder of line \} } & 733\end{array}$
else SError(4) \{ invalid character encountered \}
else

NewTokenPtr^. Char Indx := TableIndex
end
$\frac{\text { end }}{\text { se }}$
$\frac{\text { else }}{i f}$
$\frac{\text { else }}{\text { if }}$ MonadicReference
then SError(74) \{ monadic reference to dyadic operator \}
else
$\frac{\text { begin }}{\text { bew }}$ operator is dyadic \}
NewTokenPt $r^{r}$. noun $:=$ DyadOper;
NewTokenPt $r^{-}$. DOp Ind $x:=$ TableIndex
NewTokenPtr ${ }^{\wedge}$. DOp Indx $:=$ TableIndex
end
end \{ dyadicopcheck \};
procedure CheckOtherTables;
var
$\frac{\text { TableIndex: integer; }}{}$
TableIndex: integer
ChkIndex: integer;
function NextNonBlank: integer;
begin
ChkIndex : = position + 1;
while (ChkIndex < LineLength) and (APLstatement[ChkIndex] =
character[space]) do
character[space]) do
ChkIndex $:=$ ChkIndex + 1;
NextNonBLank $:=$ APLstatement[ChkIndex];
NextNonBlank $:=$ APLsta
end $\{$ nextnonblank \};
begin \{ checkothertables \}
if NextNonBLank = character[ForwardSLash]
then
beg in
TableLookUp(APLstatement[position], 16, RedTab, TableIndex);
TableLookUp(APLsta
if TableIndex $=0$
if TableIndex $=0$
then SError (72) \{ invalid reduction operator \}67677
p
$\frac{\text { begin }}{\text { new (NewVal TabLink) ; }} \begin{gathered}\text { store values }\end{gathered}$
new(NewValTabLink);
NewValTabLink^. NextValTabLink : $=01 d V a(T a b L i n k ;$
NewValTabLink
OldValTabLink : NextValTabLink : NewValTabLink;
NewValTabLink^. ForwardOrder := true;
if FunctionMode then NewValTabLink.. IntermedResult := false
$\frac{\text { if FunctionMode then NewValTabLink }}{}$ e. Interm
switch := true;
while ItsANumber do
$\frac{\text { while }}{}$ begin
begin
NumberCount := NumberCount +1 ; new(NewValues);
if switch
then beg in
beg in
switch := false;
NewValTablink ${ }^{\wedge}$. FirstValue $:=$ NewValues
end
else NewValPtr^. NextValue : $=$ NewValues;
NewValues^. RealVal := RealNumber; NewValPtr := NewValues;
NewValues^. RealVal $:=$ Real Number;
MakeNumber (Real Number, It sANumber)
end;
end;
Ne $\frac{\text { end; }}{\text { wValues }}$. NextValue $:=$ nil;
if NumberCount > 1
if Numbercount >1
if Num
then
$\frac{\frac{\text { then }}{\text { 保 }}}{\text { New }}$
677
679
679
680
681
681
682
683
$\begin{array}{ll}\text { \}; } & 711 \\ & 712 \\ & 713 \\ & 714\end{array}$
NewTokenPtr $r^{n}$.noun $:=$ SpecOper;
d \}
Lues ". NextVa
mbercount $>$
$\frac{b e s}{N}$
NewValTabLink^.dimensions $:=1$ number is a vector
NewValTabLink^. dimensions $:=1$ number is a vector
new(NewDim); NewValTabLink^. FirstDimen $:=$ NewDim;
new(NewDim); NewValTabLink^. FirstD
NewDimn
NewDim^. NextDimen := nil
el $\frac{\text { end }}{\text { Ne }}$
el $\frac{\text { end }}{\text { se }}$
$\frac{\text { else }}{\text { begin }}$
$\frac{\text { begin }}{\text { New }}$
$\frac{\text { Negin }}{\text { NewalTabLink^. .dimensions }:=0}$
NewValTabLink^.dimensions $:=0\{$ number is a scalar \};
NewValTabLink^. FirstDimen $:=$ nil
end;
NewTokenPt $r^{n}$. noun $:=$ constant;
NewTokenPtr ${ }^{\wedge}$. ValTabPtr : $=$ NewValTabLink;
Ne
end
NewTokenPtr ${ }^{\wedge}$. . . Vaun $:=$ constab
end $\frac{\text { end }}{\{ }$ trytogetanumber $\} ;$
function NameInVarTable (name: PackedString; var VarPointer:
VarTabPtrType; TestFuncPtr: PtrFuncTab): Boolean;
$\frac{\text { var }}{\text { found: Boolean; }}$
beg in
$\frac{\text { begin }}{\text { found }}:=$ false; VarPointer $:=$ OLdVarTabPtr;
found $:=$ false; VarPointer $:=$ OldVarTabPtr;
while (VarPointer <> nil) and (not found) do
$\frac{\text { beg in }}{\text { if }}$
if (NamesMatch(name, VarPointer^. VarName)) and (VarPointer.
FuncTabPtr = Test FuncPtr) \{ test for global var \}
FuncTabPtr $=$ Tes
then found $:=$ true
$\begin{aligned} & \text { then } \\ & \text { else } \\ & \text { els } \\ & \text { VarPointer }\end{aligned}:=$ VarPointer ${ }^{\wedge}$. NextVarTabPtr
end;
NameInVarTable := found;
NameInVarTable := foun
end \{ nameinvartable \};
end \{ nameinvartable \};
procedure AddNameToVarTable(name: PackedString);
begin \{ new variable name encountered \}

OldVarTabPtr := NewVarTabPtr; NewVarTabPtr". VarName := name;
OldVarTabPtr $:=$ NewVarTabPtr;
NewVarTabPtr ${ }^{\wedge}$.ValTabPtr $:=$ nil;
$\underset{\infty}{\sim} \underset{\sim}{Z}$
682
683
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690
691

```
            \frac{else}{if}
```



```
            else
```



```
            NewTokenPtr'.noun := ReductOper;
            \}\begin{array}{l}{\mathrm{ New}}\\{\mathrm{ end;}}
        position:= ChkIndex + 1;
    el的年
\frac{else}{\mathrm{ begin}}
    begin
        TableLookUp(APLstatement[position], 9, MOpTab, TableIndex);
        TableLookUp(APLstatement[position], 9,
        if TableIndex = 0 then DyadicOpCheck
            #-tse
            \frac{else}{\frac{if}{else}}\mathrm{ not MonadicReference then DyadicOpCheck}
            else
            begin { operator is monadic
            begin { operator is monadic }
                    NewTokenPtr . noun := MonadOper;
            sition; := position + 1;
            posi\frac{end;}{tion}}:=\mathrm{ position + 1;
        #
end end
end}{\mathrm{ end checkothertables };
procedure TryToGetANumber;
var
Nar
    NumberCount: integer
    RealNumber: real;
beg in
begin
    NumberCount := 0; MakeNumber(RealNumber, It sANumber);
    NumberCount := 0; MakeNumber(RealNumber,
    else
    begin { store values in value table }
            NewTokenPtr r}\mp@subsup{}{}{-}\mathrm{ .RedIndx := TableIndex;
        position := position + 1
    ItsANumber: Boolean;
```

692
693
693
694
694
695
o o o
$\begin{array}{r}697 \\ 698 \\ 699 \\ \hline\end{array}$
715
704
705
706
06
?

| 779 | if NewTokenPtr <> nil |
| :---: | :---: |
| 780 |  |
| 781 | if (NewTokenPtr ${ }^{\text {® }}$.noun $=$ FormRes) or (NewTokenPtr ${ }^{\text {® }}$. $\mathrm{noun}=$ FormArg |
| 782 | ) |
| 783 |  |
| 784 | else NewVarTabPtri. FunctabPtr : $=$ nil |
| 785 | end ( addnametovartable \}; |
| $787$ |  |
|  |  |
| 788789 | function FunctionAlreadyDefined(var NewfuncName: PackedString; var |
|  | FuncIndex: Ptrfunctab): Boolean; |
| 790 |  |
| 791 |  |
| 792 | found: Boolean; |
| 793 |  |
| 794 | begin |
| 795 found := false; FuncIndex := OldFuncTabPtr ; | found := false; FuncIndex := OldFuncTabPtr; |
| 796 while (FuncIndex <> nil) and (not found) and (NewfuncTabPtr <> |  |
| 797 | nil) do |
| 798 i¢ NamesMatch(FuncIndex* FuncName, NeuFuncName) | if NamesMatch(Func Index*.FuncName, NewFuncName) |
| 799 then found := true |  |
| 800 | else FuncIndex := FuncIndex^. NextFuncTabPtr; |
| 801 FunctionAlreadydefined := found |  |
| 802 | end ( functionalreeadydefined \}; |
| 804803 |  |
|  |  |
| 805 p | procedure MakeTokenLink; |
| 806 |  |
| 807 | begin |
| 808 new(NewTokenPtr); NewTokenPtr ${ }^{\wedge}$.NextToken := OldTokenPtr ; |  |
| 809 | SaveTokenPtr : OLdTokenPtr; OldTokenPtr : = NewTokenPtr |
| 810 end \{ maketokenlink \}; |  |
| 811 |  |
| 812 |  |
| 813 | procedure ProcessFuntionteader; |
| 814 |  |
| 816 Dummyptr: ^Functab; |  |
|  |  |
| 817 name1, name2, name3: PackedString; |  |
| 818 ItsAnIdentifier, Funcheaderror: Boolean; |  |
| 819 Arity Index: integer; |  |
| 820 begin |  |
|  |  |
| 821 begin FuncHeadError := false; FunctionMode := true; |  |
| 823 FuncStatements :=-1; |  |
| 824 | if Firstfunction |
| 825 then begin FuncStatements :=0; FirstFunction := false; end; |  |
| 826 AriTy Index $:=1$; position $:=$ position + 1; |  |
| 827 identifier(name1, ItsAnIdentifier); |  |
| 828 if not ItsAnIdentifier |  |
| 829 then |  |
| 830 begin |  |
| 831 | SError (7) \{ unrecognizable function'argument name \}; |
| 832 FunctionMode := false i exit function mode \}; |  |
| 833 FuncHeadError := true |  |
| 834 end |  |
| 834 else |  |
| 836 begin |  |
| 837 new(NewFunctabPtr); SkipSpaces; |  |
| 838 if APLstatement[position] = character[LeftArrow] |  |
|  |  |
|  |  |
| 841 NewfuncTabPtr".result := true \{ explicit result \}; |  |
| 843 ( ${ }^{\text {a }}$ identifier(name1, ItsAnIdentifier) ; |  |
|  |  |
| 8445 if not ItsAnIdentifier then |  |
|  |  |
| 847 SError (6) |  |
| 848 \{ unrecognizable name to right of explicit res \}; |  |
|  |  |
|  |  |
| 851 | $\frac{\text { end }}{\text { se }}$ Newfunctabetr ${ }^{\text {a }}$ result $:=$ false it no explicit result |
| 852 else NewFuncTabPtr ${ }^{\wedge}$.result : $=$ false \{ no explicit result \}; |  |
| 853 SkipSpaces; |  |
| $854$ |  |
| 855 | 856 begin |
| 856 |  |
| 857 identifier(name2, ItsAnIdentifier); | if not ItsAnIdentifier |
| 858 if not itsanldentifier |  |
| 859 860 $\frac{\text { then }}{\text { begin }}$ |  |
| 861 SError (7) invalid function'argument name \}; |  |
| 862 end Funcheaderror := true |  |
| 863 elst ${ }^{\text {se }}$ Arity Index $:=2$ |  |
|  |  |  |
| 864 end; ${ }^{865}$ els |  |
| 865 SkipSpaces; |  |
| 867 | if (position $<=$ LineLength) and (not Funcheaderror) |
| 868 |  |
| 869 | $\frac{\text { identifier(name }}{}$, ItsAnIdentifier) ; |
| 870 | if not ItsAnIdentifier |
| 871 | then $\frac{1}{}$ |
| 872 | beg in |
| 873 | SError (9) \{ invalid function right argument name \}; |
| 874 | FuncheadError : $=$ true |
| 875 | end |
| 876 877 | else $A r i t y$ Index $:=3$ |
| 877 | Sk $\frac{\text { end; }}{\text { ¢ }}$; ${ }^{\text {coces }}$; |880then

if (position <= LineLength) and (not FuncHeadError) then $\frac{\text { begin }}{\text { SError (3) }}$
\{ extraneous characters to right of function header
FuncheadError := true
end;
case Arity Index of
beg in
NewfuncTabPtr-.arity $:=$ niladic;
NewfuncTabPtri. FuncName := name1;
2 :
beg in
NewfuncTabPtr".arity : $=$ monadic;
NewfuncTabPtr ${ }^{\wedge}$. FuncName : $=$ name 1 ;
NewfuncTabPtr".RightArg := name2;
AddNameToVarTable(name2);
NewVarTabPtr". FuncTabPtr $:=$ NewFuncTabPtr;
$3: \frac{\text { end; }}{}$
3: beg in
NewfuncTabPtr".arity $:=$ dyadic;
NewfuncTabPtr". LeftArg := name1
NewFuncTabPtr". FuncName := name2
NewfuncTabPtr".RightArg := name3;
AddNameToVarTable(name1);
NewVarTabPtr". FuncTabPtr : = NewFuncTabPtr;
AddNameToVarTable(name3);
NewVarTabPtr ${ }^{-}$. FuncTabPtr ${ }^{\prime}:=$ NewfuncTabPtr;
end
end $T$ case ;
if FunctionAl readyDefined (NewFuncTabPtr ${ }^{-}$. FuncName, DummyPtr)
then
beg in
Error(5) \{ function already defined \};
FuncheadError := true;
end;
if FuncheadError then
beg in
dispose (NewFuncTabPtr) \{ header no good \};
FunctionMode := false \{ exit function mode \};
NewfuncTabPtr := OldfuncTabPtr;
end
end
end $\{$ processfuncheader \};
procedure DestroyStatement
$\xrightarrow[\text { Dar }]{\text { DumTokenPtr: "TokenTable; }}$
AuxSubrTabPtr: "SubrTab;
beg in
if SubrTabPtr <> nil
then
beg in
while SubrTabPtr ${ }^{\wedge}$. LastSubrPtr $<>$ nil do
begin
AuxSubrTabPtr : = SubrTabPtr
SubrTabPtr := SubrTabPtr". LastSubrPtr;
dispose(AuxSubr TabPtr)
end;
dispose(SubrTabPtr).
end;
DumTokenPtr : $=$ OldTokenPtr
while DumTokenPtr <> HoldTokenPtr do
begin
OldTokenPtr := OldTokenPtr^.NextToken; dispose(DumTokenPtr);
DumTokenPtr := OldTokenPtr
end:
NewTokenPtr := HoldTokęnPtr;
OldTokenPtr := HoldTokenPtr
\{ return pointer to end of last good line \}
end \{ destroystatement \};
procedure ReverseLinkList(var ArgPtr: TypeVa(TabPtr);
$\frac{\text { var }}{\text { hold, TemPtr: }}$ •values;

ValPtr : = ArgPtr". FirstV
while TemPtr <> nil do
begin
hold $:=$ TemPtrn.NextValue; TemPtrn.NextValue $:=$ ValPtr;
ValPtr : $=$ TemPtr; TemPtr $:=$ hold
end;
ArgPtr^. FirstValue^. NextValue : $=$ nil;
ArgPtr". FirstValue : $=$ ValPtr;
ArgPtr . FirstValue := Val
then ArgPtr ${ }^{-}$. ForwardOrder $:=$false
else ArgPtr ${ }^{\text {. }}$. ForwardOrder $:=$ true $\{$ toggle list order switch \}
end \{ reverselinklist \};
procedure parser(var TokenTabPtr: TokenPtr; var PtrToDa: TypeValTabPtr);
var

\begin{tabular}{|c|c|}
\hline 981 \& VFunchold: -vfunc \{ hold while searching \}; <br>
\hline 982 \& AuxOperTabPtr: "OperandTab; <br>
\hline 983 \& AuxSubr TabPtr: "Subrtab; <br>
\hline 984 \& AuxRParmPtr: ${ }^{\text {c FParmTab; }}$ <br>
\hline 985 \& AuxLParmPtr: ${ }^{\text {cFParmTab; }}$ <br>
\hline 986 \& ValidExp: Boolean \{ true if valid expression \}; <br>
\hline 987 \& cnt: integer; <br>
\hline 988 \& npv: integer \{ number of indices \}; <br>
\hline 989 \& assign, assign1: Boolean \{ assign.in progress \}; <br>
\hline 990 \& DoneSuccessor: Boolean; <br>
\hline 991 \& DoneParse: Boolean; <br>
\hline 992 \& <br>
\hline 993 \& <br>
\hline 994 \& procedure error(ErrorIndex: integer); <br>
\hline 995 \& <br>
\hline 996 \& var <br>
\hline 997 \& MsgCol: integer; <br>
\hline 998 \& <br>
\hline 999 \& begin <br>
\hline 1000 \& Write(' ', ErrorIndex, ' '); <br>
\hline 1001 \& for MsgCol := 1 to MessageLength do <br>
\hline 1002 \& write(ErrorMsgs[ErrorIndex, MsgCol]); <br>
\hline 1003 \& writeln; goto 100 i return to scanner \}; <br>
\hline 1004 \& end ( error ${ }^{\text {T }}$; <br>
\hline 1005 \& <br>
\hline 1006 \& <br>
\hline 1007 \& procedure retease; <br>
\hline 1008 \& <br>
\hline 1009 \& begin $\{$ releaseopertab \} <br>
\hline 1010 \& OperTabPtr : = PtrLastOper; <br>
\hline 1011 \& while OperTabPtr ${ }^{\text {a }}$.LastOper <> nil do <br>
\hline 1012 \& beg in <br>
\hline 1013 \& AuxOperTabPtr := OperTabPtr; <br>
\hline 1014 \& OperTabPtr : = OperTabPtr ${ }^{\text {-.LastOper }}$; dispose(AuxOperTabPtr); <br>
\hline 1015 \& end; <br>
\hline 1016 \& end Treleaseopertab \}; <br>
\hline 1017 \& <br>
\hline 1018 \& <br>
\hline 1019 \& procedure expression(var ValidExp: Boolean); <br>
\hline 1020 \& forward; <br>
\hline 1021 \& <br>
\hline 1022 \& <br>
\hline 1023 \& procedure ReturntocallingSubr; <br>
\hline 1024 \& <br>
\hline 1025 \& var <br>
\hline 1026 \& NamePtr: `VarTab; <br>
\hline 1027 \& <br>
\hline 1028 \& begin \{ returntocallingsubr \} <br>
\hline 1029 \& if SubrTabPtr ${ }^{\text {n }}$. CalledSubr ${ }^{\text {², result }}$ <br>
\hline 1030 \& then <br>
\hline 1031 \& begin \{ place explicit result in opertab \} <br>
\hline 1032 \&  <br>
\hline 1033 \& NamePtr, SubrTabPtr ${ }^{\circ}$. CalledSubr) <br>
\hline 1034 \& then error(11) ( 'symbol not found' <br>
\hline 1035 \& else <br>
\hline 1036 \& begin <br>
\hline 1037 \& AuxOperTabPtr : = OperTabPtr; new(OperTabPtr); <br>
\hline 1038 \& OperTabPtr ${ }^{\text {n }}$-LastOper $:=$ AuxOperTabPtr; <br>
\hline 1039 \& PtrLastoper : = OpertabPtr; <br>
\hline 1040 \& OperTabPtr ${ }^{\text {n }}$. OperPtr $\mathrm{:}$ : NamePt $r^{\sim}$.ValTabPtr; <br>
\hline 1041 \& end; <br>
\hline 1042 \& \{ $\frac{\text { end }}{\text { return to calling function }}$, <br>
\hline 1043 \&  <br>
\hline 1044 \& VFuncPtr $:=$ Subr TabPtr ${ }^{\text {a }}$. StatemCallingSubr; <br>
\hline 1045 \& TokenTabPtr $:=$ SubrTabPt $r^{-}$. TokencallingSubr ${ }^{\text {n }}$. NextToken; <br>
\hline 1046 \& if SubrTabPtr".CalledSubr^.arity <> niladic <br>
\hline 1047 \& then <br>
\hline 1048 \& begin \{ monadic or dyadic \} <br>
\hline 1049 \& AuxRParmPtr : = RParmPtr; RParmPtr := RParmPtr ${ }^{\wedge}$.LastParm; <br>

\hline 1050 \& | dispose(AuxRParmPtr); |
| :--- |
| if SubrTabPtr".CalledSubr^.arity $=$ dyadic then | <br>

\hline 1051 \& If Subr TabPtr .CalledSubr .arity = dyadic then <br>
\hline 1052 \& AuxLParmPtr : = LParmPtr; <br>
\hline 1054 \& LParmPtr := LParmPtr ${ }^{\text {- }}$.LastParm; dispose(AuxLParmPtr); <br>
\hline 1055 \& end; <br>
\hline 1056 \& Aux SubrtabPtr : $=$ SubrTabPtr ; <br>
\hline 1057 \& SubrtabPtr : $=$ SubrtabPt $r^{-}$.LastSubrPtr ; dispose(AuxSubrTabPtr) ; <br>
\hline 1058 \&  <br>
\hline 1059 \& end \{ returntocallingsubr \}; <br>
\hline 1060 \& <br>
\hline 1061 \& function SpecSymbol(sym: integer): Boolean; <br>
\hline 062 \& function Specsymbolssm. integer): Boolean, <br>
\hline 063 \& var <br>
\hline 1064 \& ValidSym: Boolean; <br>
\hline 066 \& <br>

\hline 1067 \& | begin specsymbol |
| :--- |
| ValidSym := false; | <br>

\hline 1068 \& if TokenTabPt ${ }^{\wedge}$. ${ }^{\text {noun }}=$ SpecOper <br>
\hline 1069 \& then <br>
\hline 1070 \& <br>
\hline 1071
1072 \& begin <br>
\hline 1072 \& hold := TokenTabPtr; <br>
\hline 1073 \& TokenTabPtr : = TokenTabPtr ${ }^{\wedge}$.NextToken; ValidSym := true; <br>
\hline 1074 \& end; <br>
\hline 1075
1076 \& SpecSymbol := ValidSym; <br>
\hline 1077 \& end \{ specsymbol \}; <br>
\hline 1078 \& <br>
\hline 1079
1080 \& procedure Callsubr; <br>
\hline
\end{tabular}

procedure error(ErrorIndex: integer);
$\frac{\text { var }}{\text { MsgCol: integer; }}$
begin
write(' ', ErrorIndex, ' ');
for MsgCol $:=1$ to MessageLength do
write(ErrorMsgs[ErrorIndex, Msg $\frac{\text { Col }}{}$ );
writeln; goto 100 \{ return to scanner \};
writeln; goto
procedure retease
begin \{ releaseopertab \}
OperTabPtr : = PtrLastOper:
while OperTabPtr. LastOper <> nil do
beg in
AuxOperTabPtr : $=$ OperTabPtr:
OperTabPtr $:=$ OperTabPtr ${ }^{-}$. LastOper; dispose(AuxOperTabPtr);
end;
procedure expression(var ValidExp: Boolean);
forward;
procedure ReturnToCallingSubr:
NamePtr: "VarTab;
begin \{ returntocallingsubr \}
if SubrTabPtrn. CalledSubr".result
then
begin \{ place explicit result in opertab \}
if not Name InVarTable (SubrTabPtr". CalledSubr^. Resul tName,
NamePtr, SubrTabPtr". CalledSubr)
then error(11) \{ 'symbol not found' \}
else
Aegin $\quad$ OperTabPtr := OperTabPtr; new(OperTabPtr).
OperTabPtr". LastOper $:=$ AuxOperTabPtr:
OperTabPtr . LastOper $:=$ Aux
PtrLastOper $:=$ OperTabPtr;
OperTabPtr ${ }^{\wedge}$. OperPtr $:=$ NamePtr ${ }^{\wedge}$. ValTabPtr.
end;
end;
return to calling function \}
VFuncPtr := SubrTabPtr". StatemCallingSubr;
TokenTabPtr : = SubrTabPtr ${ }^{-}$. TokenCallingSubr ${ }^{\wedge}$. NextToken;
if SubrTabPtr^. CalledSubr^.arity <> nitadic
begin \{ monadic or dyadic \}
AuxRParmPtr $:=$ RParmPtr; RParmPtr $:=$ RParmPtrn.LastParm;
AuxRParmPtr := RParmPt
dispose(AuxRParmPtr);
if SubrTabPtr $r^{-}$. CalledSubr - .arity $=$dyadic then
begin \{ dyadic only \}
AuxLParmPtr: $=$ LParmPtr;
end;
end;
LParmPtr := LParmPtr". LastParm;
dispose(AuxLParmPtr)
AuxSubrTabPtr := SubrTabPtr;
SubrTabPtr := SubrTabPtr". LastSubrPtr; dispose(AuxSubrTabPtr);
end \{ returntocallingsubr \};
function SpecSymbol(sym: integer): Boolean;
$\frac{\operatorname{var}}{\mathrm{V}}$
ValidSym: Boolean;
begin \{ specsymbol \}
VatidSym := false;
if TokenTabPtr ${ }^{-}$.noun $=$SpecOper
then
if TokenTabPtr^. Char Indx $=$ sym then
$\frac{\text { begin }}{\text { hold }}:=$ TokenTabPtr;
TokenTabPtr := TokenTabPtr^.NextToken; ValidSym := true;
end;
Symbol := ValidSym;
AuxOperTabPtr: "OperandTab;
AuxSubrTabPtr: "SubrTab;
AuxRParmPtr: *FParmTab;
AuxRParmPtr: "FParmTab;
AuxLParmPtr: "FParmTab;
Validexp: Boolean $\{$ true if valid expression \};
cnt: integer:
npv: integer \{ number of indices \};
assign, assign1: Boolean \{ assign.in progress \};
DoneSuccessor: Boolean
DoneParse: Boolean
rocedure CallSubr;

```
3 Var PtrToVarTab: ^VarTab;
3 Var PtrToVarTab: `VarTab;
1085 begin \(\{\) callsubr \(\}\)
    if SubrTabPtr \({ }^{\wedge}\). CalledSubr \({ }^{\wedge}\).arity <> niladic
    then
        begin
```




```
1091 then error(32); \(\quad\) if PtrToVarTab*. FuncTabPtr <> SubrTabPtr^. CalledSubr
        then error(32) \{ program logic error, variable name of \};
        \{ function argument not found in symbol table \}
            AuxRParmPtr := RParmPtr; new(RParmPtr);
            RParmPtr \({ }^{-}\).LastParm :=AuxRParmPtr;
        PtrToVarTab \({ }^{\wedge}\). DeferedValTabPtr \(:=\) RParmPtr;
        if SubrTabPtr \({ }^{\wedge}\). CalledSubr \({ }^{\wedge}\).arity \(=\) dyadic
        then
            begin \{ if dyadic \}
                    if not Name In VarTable (SubrTabPtr". CalledSubr \({ }^{\text {. }}\).LeftArg,
                    - PtrToVarTab, SubrTabPtr". CalledSubr)
                    then error(33);
                    if PtrToVarTab". FuncTabPtr <> SubrTabPtr^. CalledSubr
                    then error(33) \{ same as error(32) \};
                    AuxLParmPtr \(:=\) LParmPtr; new(LParmPtr);
LParmPtr \({ }^{2}\).LastParm \(:=\) AuxLParmPtr;
                    AuxLParmPtr \(:=\) LParmPtr; new(LParmPtr);
LParmPtr
                    LParmPtr". LastParm \(:=\) AuxLParmPtr;
PtrToVarTab". DeferedValTabPtr : LParmPtr;
                    LParmPtr". Ptrval : = OperTabPtr". OperPtr;
                    AuxOperTabPtr := OperTabPtr;
                    OperTabPtr : = OperTabPtr". LastOper;
                    dispose(AuxOperTabPtr); PtrLastÓper : = OperTabPtr;
                end;
            RParmPtr^.PtrVal : = OperTabPtr^. OperPtr;
            AuxOperTabPtr: = OperTabPtr;
            AuxOperTabPtr : = OperTabPtr;
OperTabPtr : \(=\) OperTabPtr \({ }^{-}\).Last Oper; dispose(AuxOperTabPtr);
            PtrLastOper := OperTabPtr;
        To \(\overline{k e n T a b P t r}:=\) SubrTabPtr^. CalledSubr ^. FirstStatement \({ }^{\wedge}\). NextStmnt;
    VFuncPtr := SubrTabPtr".CalledSubr^.FirstStatement;
    end \{ callsubr \};
function functCall: Boolean;
    \(\frac{\text { var }}{P}\)
        PtrToFuncTab: "FuncTab;
        NameOfFunc: PackedString;
        ValidFn: Boolean;
    begin \{ functcall \}
    Validfn := false;
        if TokenTabPtr". noun \(=\) GlobVar
        then
            begin
            NameOfFunc := TokenTabPtr \({ }^{\wedge}\). VarTabPt \(r^{\wedge}\). VarName;
                if FunctionAlreadyDefined(NameOf Func, PtrToFunc Tab)
                then
                    AuxSubrTabPtr : = SubrTabPtr; new(SubrTabPtr);
                    SubrTabPtr \({ }^{\wedge}\). LastSubrPtr : = AuxSubrTabPtr;
                    SubrTabPtr \({ }^{- \text {. CalledSubr }:=P t r T o F u n c T a b ; ~}\)
                    SubrTabPtr \({ }^{\wedge}\). CalledSubr \(:=\) PtrToFuncTab;
SubrTabPtr \({ }^{\wedge}\). TokenCallingSubr \(:=\) TokenTabPtr;
                    SubrTabPtr \({ }^{-}\). StatemCallingSubr \(:=\)VFuncPtr;
                    hold := TokenTabPtr;
                    TokenTabPtr : = TokenTabPtr \({ }^{\wedge}\).NextToken; ValidFn := true;
        end; \({ }^{\text {end; }}\)
    FunctCall := ValidFn;
    end \{ functcall \};
    procedure NunWrite(RealNo: real);
    \(\frac{\mathrm{var}}{\mathrm{pr}}\)
        \(\frac{\text { var }}{\text { prefix, root: integer; }}\)
        SigDig, ColCnt: integer;
    begin \{ output a number \}
        if Real No \(>=0.0\)
        then write(1 ', RealNo: 12: 2) \{ output positive number \}
        \(\frac{\text { else }}{}\)
        begin \(\{\) output negative number \}
            Real No : \(=-1.0\) *RealNo;
            SigDig \(:=\operatorname{trunc}\left(\left(\ln \left(\operatorname{Rea}\left(N_{0}\right)\right) /(\ln (10.0))\right)\right.\);
            for ColCnt :=1 to(7-SigDig) do write(' ');
            \(\frac{\text { if }}{}\) character[negative] \(<6000\)
            then write(chr(character[negative]))
            else
                    beg in
```



```
                    prefix \(:=\) character[negative] div 100;
root \(:=\) character[negative] \(-(100 \star\) prefix \()\);
                    write(chr(prefix), chr(root));
                end;
            SigDig \(:=\) SigDig +5 ; write(RealNo: SigDig: 2 );
        end
procedure OutPutVal;
1083
1086
1087
1088
1093
1094
    end \(?\)
        if PtrToVarTab*. FuncTabPtr <> SubrTabPtr .CalledSubr
            then
        end;
                    beg in
                    AuxSubrTabPtr := SubrTabPtr; new(SubrTabPtr)
        end;
1181
```1163
1164
1165
```

    SigDig, ColCnt: integer,
        numwrite
            te \};
    ```
\begin{tabular}{|c|c|}
\hline 1182 & var \\
\hline 1183 & cnt: integer; \\
\hline 1184 & AuxValuesptr: "values; \\
\hline 1185 & DimHold, dimen1, dimen2, dimen3: integer; \\
\hline 1186 & OutCnt1, Out Cnt2, OutCnt3: integer; \\
\hline 1187 & idimens: integer; \\
\hline 1188 & \\
\hline 1189 & begin \{ outputval \\
\hline 1190 & cnt := 0; writeln; writeln; \\
\hline 1191 & if not OperTabPtr \({ }^{\text {n }}\).OperPtr \({ }^{\text {n }}\).ForwardOrder \\
\hline 1192 & then ReverseLinkList (OperTabPtr \({ }^{\wedge}\). OperPtr) ; \\
\hline 1193 & AuxValuesPtr : = OperTabPtr \({ }^{\text {a }}\). OperPtr \({ }^{\text {a }}\).FirstValue; \\
\hline 1194 & idimens : \(=\) OperTabPtr \({ }^{\text {® }}\).OperPtr \({ }^{\wedge}\).dimensions; \\
\hline 1195 & if not (idimens in [0..3]) \\
\hline 1196 & then \\
\hline 1197 & begin \\
\hline 1198 & for colcnt \(:=1\) to MessageLength do \\
\hline 1199 & write(ErrorMsgs[60, ColCnt]); \\
\hline 1200 & writeln; \\
\hline 1201 & end \\
\hline 1202 & else \\
\hline 1203 & \(\underline{\text { if }}\) AuxValuesPtr \(=\) nil \\
\hline 1204 & then \\
\hline 1205 & begin \\
\hline 1206 & for colcnt : \(=1\) to MessageLength do \\
\hline 1207 & write(ErrorMsgs[61, colCnt]); \\
\hline 1208 & writeln; \\
\hline 1209 & end \\
\hline 1210 & else \\
\hline 1211 & if idimens \(=0\) \\
\hline 1212 & then begin NunWrite(AuxValuesPtr \({ }^{\text {a }}\).RealVal); writeln; end \\
\hline 1213 & else \\
\hline 1214 & begin \\
\hline \[
\begin{aligned}
& 1215 \\
& 1216
\end{aligned}
\] & dimen1 \(:=\) OperTabPtr^. OperPtr^. FirstDimen^. \({ }^{\text {dimenlength; }}\) if idimens >= 2 \\
\hline 1217 & then \\
\hline 1218 & dimen2 : = OperTabPtr \({ }^{\text {® }}\).OperPtr \({ }^{\text {® }}\).FirstDimen \({ }^{\text {a }}\).NextDimen \\
\hline 1219 & *.dimenlength \\
\hline 1220 & else dimen2 : \(=1\); \\
\hline 1221 & if idimens \(=3\) \\
\hline 1222 & then \\
\hline 1223 &  \\
\hline 1224 & *.NextDimen \({ }^{\text {- }}\) dimenlength \\
\hline 1225 & else dimen 3 : \(=1\); \\
\hline 1226 & if idimens \(=3\) then \\
\hline 1227 & begin \(\{\) rotate dimensions \\
\hline 1228 & DimHold := dimen1; dimen \(:=\) dimen2; \\
\hline 1229 & dimen2 := dimen3; dimen3 := DimHold; \\
\hline 1230 & end; \\
\hline 1231 & for OutCnt3 := 1 to dimen3 do \\
\hline 1232 & begin \\
\hline 1233 & for OutCnt2 := 1 to dimen \({ }^{\text {do }}\) \\
\hline 1234 & begin \\
\hline 1235 & for OutCnt1 \(:=1\) to dimen2 do \\
\hline 1236 & begin \\
\hline 1237 & cnt \(:=\mathrm{cnt}+1\); \\
\hline 1238 & if ( ( cont - 1) \(\bmod 5)=0\) ) and (cnt <> 1) \\
\hline 1239 & then begin writetn; write(' '); end; \\
\hline 1240 & NunWrite(AuxValuesPtr \({ }^{\text {- }}\).RealVal); \\
\hline 1241 & AuxValuesPtr : = AuxValuesPtr \({ }^{\text {a }}\).NextValue; \\
\hline 1242 & end; \\
\hline 1243 & if idimens >= 2 \\
\hline 1244 & then begin writeln; cnt : \(=0\); end; \\
\hline 1245 & end; \\
\hline 1246 & writeln; writeln; \\
\hline 1247 & end; \\
\hline 1248 & \{ writeln; \} \\
\hline 1249 & end; \\
\hline 1250 & end \{ outputval \}; \\
\hline 1251 & \\
\hline 1252 & \\
\hline 1253 & function variable: Boolean; \\
\hline 1254 & \\
\hline 1255 & var \\
\hline 1256 & globOrDummy: Boolean \{ gord \}; \\
\hline 1257 & Passedadj: 'Vartab \{ k \}; \\
\hline 1258 & rarg: Boolean \{ rd \}; \\
\hline 1259 & ParmPtr: "Valtab ( pt \}; \\
\hline 1260 & VatidVar: Boolean; \\
\hline 1261 & ValidIndex: Boolean; \\
\hline 1262 & \\
\hline 1263 & \\
\hline 1264 & procedure InputVal; \\
\hline 1265 & \\
\hline 1266 & var \\
\hline 1267 & AuxPtrToda: "Valtab; \\
\hline 1268 & AuxValuesPtr: "values; \\
\hline 1269 & Aux2valuesPtr: "values; \\
\hline 1270 & Realv: real; \\
\hline 1271 & boolv: Boolean; \\
\hline 1272 & contr, cnt: integer; \\
\hline 1273 & AuxDimenFoptr: "DimenInfo; \\
\hline 1274 & \\
\hline 1275 & begin \{ inputval \} \\
\hline 1276 & Cnt \(:=0\); position \(:=1 ;\) AuxPtrToDa \(:=\) PtrToda; \\
\hline 1277 & new(PtrToDa); AuxPtrToda^.NextValTablink := PtrToda; \\
\hline 1278 & AuxOperTabPtr : = OperTabPtr; new(OperTabPtr); \\
\hline 1279 & PtrLastoper : = OperTabPtr; \\
\hline 1280 & OperTabptr \({ }^{\text {n }}\).Lastoper : \(=\) AuxOperTabPtr; \\
\hline 1281 & OperTabPtr \({ }^{\text {a }}\). OperPtr \(:=\) PtrToDa; new(Aux2ValuesPtr);
PtrToDa*.FirstValue \(:=\) Aux2ValuesPtr; \\
\hline
\end{tabular}


Software Tools
\begin{tabular}{|c|c|}
\hline 1384
1385 & var \\
\hline 1386 & AuxPtrtoda: "Valtab; \\
\hline 1387 & PtrToValues, AuxValuesPtr: `values; \\
\hline 1388 & \\
\hline 1389 & begin \{ stackpointers \} \\
\hline 1390 & ; \\
\hline 1391 & if npv \(=0\) \\
\hline 1392 & then \\
\hline 1393 & begin \\
\hline 1394 & AuxOperTabPtr : = OperTabPtr; new(OperTabPtr); \\
\hline 1395 & OperTabPtr \({ }^{-}\). LastOper \(:=\)Aux OperTabPtr; \\
\hline 1396 & OperTabPtr^. OperPtr : \(=\) ParmPtr; \\
\hline 1397 & PtrLastOper := OperTabPtr \\
\hline 1398 & end \\
\hline 1399 & else \\
\hline 1400 & beg in \\
\hline 1401 & Auxptrtoda := PtrToda; new(PtrToda); \\
\hline 1402 & PtrTodan. NextValTablink := AuxPtrToda; \\
\hline 1403 & PtrToda . IntermedResult : = true; \\
\hline 1404 & PtrToda^.dimensions := 0; PtrToDa^.FirstDimen := nil; \\
\hline 1405 & PtrTodan. Forwardorder := true; new(AuxValuesPtr); \\
\hline 1406 & PtrToda^.FirstValue : = AuxValuesPtr; \\
\hline 1407 & GetArrayPosition(PtrToValues) ; \\
\hline 1408 & PtrTodan.FirstValuen. Realval : = PtrToValues^.RealVal; \\
\hline 1409 & PtrToda^.FirstValuen. NextValue : = nil; \\
\hline 1410 & AuxOperTabPtr : = OperTabPtr; new(OperTabPtr); \\
\hline 1411 & OperTabPtr \({ }^{\wedge}\)-LastOper \(:=\) AuxOperTabPtr ; \\
\hline 1412 & OperTabPtr \({ }^{\text {n }}\). OperPtr \(:=\) PtrToda; \\
\hline 1413 & PtrLastOper : = OperTabPtr; \\
\hline 1414 & end; \\
\hline 15 & end (stackpointers \}; \\
\hline 1416 & \\
\hline 1417 & \\
\hline 1418 & function SimpleVariable: Boolean; \\
\hline 1419 & \\
\hline 1420 & \(\frac{\mathrm{Var}}{\text { ValidSv: }}\) Boolean; \\
\hline & \\
\hline 1422 & begin \{ simplevariable \} \\
\hline 1424 & Validsv := false; rarg := false; globOrDummy := false; \\
\hline 1425 & if assign \\
\hline 1426 & then \\
\hline 1427 & begin \\
\hline 1428 &  \\
\hline 1429 & GlobVar) \\
\hline 1430 & then \\
\hline 1431 & begin \\
\hline 1432 & glob0rDummy : \(=\) true; \\
\hline 1433 & PassedAdj : = TokenTabPtr \({ }^{\text { }}\).VarTabPtr ; \\
\hline 1434 & hold : = TokenTabPtr; \\
\hline 1435 & TokenTabPtr : = TokenTabPtr \({ }^{\text {- }}\).NextToken; \\
\hline 1436 & ValidSv := true \\
\hline 1437 & end \\
\hline 1438 & else \\
\hline 1439 & if TokenTabPtr \({ }^{\wedge}\). .noun \(=\) FormArg \\
\hline 1440 & then \\
\hline 1441 & begin \\
\hline 1442 &  \\
\hline 1443 &  \\
\hline 1444 & then rarg : \(=\) true; \\
\hline 1445 & PassedAdj : \(=\) TokenTabPtr \({ }^{\wedge}\).VarTabPtr \\
\hline 1446 & end \\
\hline 1447 & end \\
\hline 1448 & else \\
\hline 1449 & begin \\
\hline 1450 &  \\
\hline 1451 & Globvar) \\
\hline 1452 & then \\
\hline 1453 & begin \\
\hline 1454 & ParmPtr : = TokenTabPtr \({ }^{\text {n }}\).VarTabPtr \({ }^{\text {n }}\). ValTabPtr ; \\
\hline 1455 & if ParmPtr <> nil then \\
\hline 1456 & begin \\
\hline 1457 & hold : = TokenTabPtr; \\
\hline 1458 & TokenTabPtr : \(=\) TokenTabPtr \({ }^{\text {². }}\). NextToken; \\
\hline 1459 & ValidSv := true \\
\hline 1460 & end \\
\hline 1461 & end \\
\hline 1462 & else \\
\hline 1463 & begin \\
\hline 1464 & if TokenTabPtr \({ }^{-}\).noun \(=\)FormArg \\
\hline 1465 & then \\
\hline 1466 &  \\
\hline 1467
1468 & - LeftArg, TokenTabPt \(r^{\wedge}\). VarTabPtr \({ }^{\wedge}\). VarName) \\
\hline 1469 & then ParmPtr \(:=\) LParmPtr \({ }^{n}\). Ptrval \\
\hline 1470 & else Parmptr \(:=\) RParmPtr \({ }^{\text {a }}\).PtrVal; \\
\hline 1471 & hold \(:=\) TokenTabPtr; \\
\hline 1472 & TokenTabPtr \(:=\) TokenTabPtr \({ }^{\text {n }}\). NextToken; \\
\hline 1473 & ValidSv := true; \\
\hline 1474 & end; \\
\hline 1475 & end; \\
\hline 1476 &  \\
\hline 1477 & SimpleVariable := ValidSv; \\
\hline 1478 & end \{ simple variable \}; \\
\hline 1479 & \\
\hline 1480 & procedure index(var ValidI: Boolean); \\
\hline 1481 & procedure index (var Vatidi: Boolean); \\
\hline 1482 & \\
\hline 1483
1484
1485 & \(\frac{\text { Var }}{}\) Valide1, Valide2: Boolean; \\
\hline & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|}
\hline 1486 & begin \{ index \} \\
\hline 1487 & ValidI := false; expression(ValidE1); \\
\hline 1488 & if Validel \\
\hline 1489 & then \\
\hline 1490 & begin \\
\hline 1491 & npv := 1 \{ no. of index expressions \}; \\
\hline 1492 & while SpecSymbol(XSemicolon) do \\
\hline 1493 & begin \\
\hline 1494 & npv := npv + 1; expression(ValidE2); \\
\hline 1495 & if not Valide2 then error(39); \\
\hline 1496 & ( 'invalid index expression' \\
\hline 1497 & end; \\
\hline 1498 & Validi := true; \\
\hline 1499 & end; \\
\hline 1500 & end \{index \}; \\
\hline 1501 & \\
\hline 1502 & \\
\hline 1503 & begin \{ variable \} \\
\hline 1504 & Validvar := false; npv := 0; \\
\hline 1505 & if not assign \\
\hline 1506 & then \\
\hline 1507 & if SpecSymbol (XQuadSym) \\
\hline 1508 & then begin Inputval; ValidVar := true end \\
\hline 1509 & else \\
\hline 1510 & beg in \\
\hline 1511 & if SpecSymbol (XRightBracket) \\
\hline 1512 & then \\
\hline 1513 & beg in \\
\hline 1514 & index(ValidIndex); \\
\hline 1515 & if (not Validindex) or (not SpecSymbol(XLeftBracket)) \\
\hline 1516 & then error(34) \{ invalid index expression \}; \\
\hline 1517 & end; \\
\hline 1518 & if SimpleVariable \\
\hline 1519 & then begin StackPointers; ValidVar := true end \\
\hline 1520 & end \\
\hline 1521 & else \\
\hline 1522 & if SpecSymbol(XQuadSym) \\
\hline 1523 & then begin OutPutVal; ValidVar := true end \\
\hline 1524 & else \\
\hline 1525 & begin \\
\hline 1526 & if SpecSymbol(XRightBracket) \\
\hline 1527 & then \\
\hline 1528 & begin \\
\hline 1529 & index(ValidIndex); \\
\hline 1530 & if (not Validindex) or (not SpecSymbol(XLeftBracket)) \\
\hline 1531 & then error(34) \{ invalid index expression \}; \\
\hline 1532 & end; \\
\hline 1533 & if SimpleVariable \\
\hline 1534 & then begin LinkResults; ValidVar := true; end; \\
\hline 1535 & end; \\
\hline 1536 & variable : = ValidVar; \\
\hline 1537 & end \(\{\) variable \}; \\
\hline 1538 & \\
\hline 1539 & \\
\hline 1540 & procedure primary(var valid: Boolean) \{ recursive entry \}; \\
\hline 1542 & \\
\hline 1543 & Validx: Boolean; \\
\hline 1544 & assign: Boolean; \\
\hline 1545 & \\
\hline 1546 & \\
\hline 1547 & function vector: Boolean; \\
\hline 1548 & \\
\hline 1549 & var \\
\hline 1550 & vec: Boolean; \\
\hline 1551 & \\
\hline 1552 & begin \(\{\) vector \} \\
\hline 1553 & vec : \(=\) false; \\
\hline 1554 & if TokenTabPtr \({ }^{\wedge}\). noun \(=\) constant \\
\hline 1555 & then \\
\hline 1556 & beg in \\
\hline 1557 & AuxOperTabPtr : = OperTabPtr; new(OperTabPtr); \\
\hline 1558 & PtrLastoper : = OperTabPtr; \\
\hline 1559 & OperTabPtr \({ }^{\text {a }}\).LastOper \(:=\) AuxOperTabPtr ; \\
\hline 1560 & OperTabPtr \({ }^{\text {a }}\). OperPtr \(:=\) TokenTabPtr \({ }^{\wedge}\).ValTabPtr; \\
\hline 1561 & hold : = TokenTabPtr; \\
\hline 1562 & TokenTabPtr : \(=\) TokenTabPtr \({ }^{\text {- }}\).NextToken; vec \(:=\) true \\
\hline 1563 & vend; \(\frac{\text { end }}{\text { ctor }}\) : \(=\) vec . \\
\hline 1564
1565 & vector := vec; end \{ vector \}; \\
\hline 1565 & end \{ vector \}; \\
\hline 1566 & \\
\hline 1567 & begin \{ primary \} \\
\hline 1569 & vatid : = true; \\
\hline 1570 & if not vector \\
\hline 1571 & then \\
\hline 1572 & beg in \\
\hline 1573 & assign := false; \\
\hline 1574 & if not variable \\
\hline 1575 & \(\frac{\text { then }}{\text { if }}\) SpecSymbol (XRightPar) \\
\hline 1576 & if SpecSymbol(XRightPar) \\
\hline 1577 & \(\frac{\text { then }}{\text { beg in }}\) \\
\hline 1578 & \(\frac{\text { begin }}{\text { expression }}\) (Validx) \\
\hline 1579 & \begin{tabular}{l}
expression(ValidX); \\
if not ValidX
\end{tabular} \\
\hline 1580
1581 & then \({ }^{\text {er }}\) eror(14) \{ 'non-valid exp within parens' \\
\hline 1582 & \(\frac{\overline{e l s e}}{\text { iff }}\) \\
\hline 1583 & \(\frac{\text { if }}{\text { then }}\) not SpecSymbol(XLeftPar) \\
\hline 1584 & \(\frac{\text { then }}{\text { error (15) }}\) \\
\hline 1585 & I 'right paren not balanced with left paren' \\
\hline & else valid := true \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline 1588 & end & 1689 & end ( dop \}; \\
\hline 1589 & else & 1690 & - \({ }^{\text {a }}\) \\
\hline 1590 & if not functcall then valid := false & 1691 & \\
\hline 1591 & else begin Callsubr; primary(valid); end; & 1692 & function ItsBoolean(test: real): Boolean; \\
\hline 1592 & end; & 1693 & \\
\hline 1593 & end Tprimary ! & 1694 & beg in \\
\hline 1594
1595 & & 1695 & if (test \(=1.0\) ) or (test \(=0.0\) ) then ItsBoolean : \(=\) true \\
\hline 1596 p & procedure expression \{ recursive \}; & 1696 & else ItsBoolean \(:=\) fa \\
\hline 597 & & 1697 & end ( itsboolean ); \\
\hline 598 & var & 1698 & \\
\hline 1599 & Done Exp, ValidPri, Validfunc, ValidAssn: Boolean; & 1699 & procedure DyadComp(var SFloat: real; value: real; code: integer); \\
\hline 1600 & code: integer; & 170 & compute result of dyadic operation \} \\
\hline 1601 & & 1702 & \\
\hline 1602 & & 1703 & beg in \\
\hline 1603 & procedure assignment (uar valida: Boolean); & 1704 & case code of \\
\hline 1604 & & 1705 & \{ left codes - reduction ops , right codes - dyadic ops \} \\
\hline 1605 & begin assignment & 1706 & 2, 52: SFloat := value + SFloat \{ addition \}; \\
\hline 1606 & valida := false; & 1707 & 3, 53: Sfloat := value - Sfloat \{ subtraction \}; \\
\hline 1608 & \(\frac{17}{\text { then }}\) SpecSymbol(XLeftarrow) & 1708 & 4, 54: SFloat := value * SFloat \{ multiplication \}; \\
\hline 1609 & begin & 1709 & 5, 55: \\
\hline 1610 & assign := true; assign \(:=\) true; & 1711 & then error(20) ( attempted division by zero \} \\
\hline 1611 & if variable then valida \(:=\) true & 1712 & else SFloat : = value / SFloat \{ division ); \\
\hline 1612 & else error(8) T result of an assn not a valid variable \}; & 1713 & 6, 56: \\
\hline 1613 & valida := true; assign := false; & 1714 & if value > 0.0 \\
\hline 1614 & end; & 1715 & then \\
\hline 1615 & end \(\{\) assignment \}; & 1716 & SFloat : \(=\exp (S F l o a t\) * \(\ln (\) vatue ) ) \\
\hline 1616 & & 1717 & \{ number raised to a power \} \\
\hline 1617
1618 & & 1718 & else SFloat := 1.0.1 (exp(SFloat * \(\ln (\) abs \((\) value ) ) ) ) \\
\hline 1618 & function mop: Boolean; & 1719 & 21, 71: \\
\hline 1619 & & 1720 & \(\underline{i f}\) value \(=\) SFloat \{ equality \} then SFloat \(:=1.0\) \\
\hline 1620 & \(\frac{\text { Var }}{\text { ValidM: }}\) Boolean; & 1721 & else \(\frac{\text { els }}{72}\) SFloat : \(=0.0\); \\
\hline 1622 & & 1722 & 22, if value <> Sfloat \{ inequality \} then SFloat := 1.0 \\
\hline 1623 & begin \{ mop \} & 1723 & else SFloat :=0.0; \\
\hline 1624 & ValidM : \(=\) false; & 1725 & 23, 73: \\
\hline 1625 & if (TokenTabPtr \({ }^{\text {² }}\).noun \(=\) MonadOper \()\) or (TokenTabPtr \({ }^{\wedge}\).noun \(=\) & 17725 & if value < Sfloat \{ less than \} then Sfloat := 1.0 \\
\hline 1626 & Reductoper) & 1727 & else SFloat : \(=0.0\); \\
\hline 1627 & then & 1728 & 24, 74: \\
\hline 1628 & \(\frac{\text { begin }}{\text { if }}\) & 1729 & if value <= Sfloat \{ less than or equal to \} \\
\hline 1629 &  & 1730 & then SFloat : \(=1.0\) \\
\hline 1630 &  & 1731 & else SFloat : \(=0.0\); \\
\hline 1631
1632 &  & 1732 & 25, 75 : \\
\hline 1632
1633 & TokenTabPtr : = TokenTabPtr \({ }^{\wedge}\).NextToken; ValidM : \(=\) true; & 1733 & if value \(>=\) SFloat \(\{\) greater than or equal to \} \\
\hline 1634 & end; & 1734 &  \\
\hline 1635 & mop := ValidM; & 1735 & \[
26,76:
\] \\
\hline 1636 & end \{ mop \}; & 1736 & if value > SFloat \{ greater than \} then Sfloat :=1.0 \\
\hline 1637 & & 1738 & else SFloat : \(=0.0\); \\
\hline 1638 & & 1739 & 27, 77: \\
\hline 1639 & function dop: Boolean; & 1740 & if (ItsBoolean(value)) and (ItsBoolean(SFloat)) \\
\hline 1640 & & 1741 & then \\
\hline 1641 & \(\frac{\mathrm{var}}{\text { Validd: }}\) Boolean; & 1742 & \(\frac{\text { if }}{}\) (value \(=1.0\) ) and (SFloat \(=1.0\) ) \(\{\) and \\
\hline 1642 & Valido: Boolean; & 1743 & then SFloat \(:=1.0\) \\
\hline 1643
1644 & & 1744 & else SFloat \(:=0.0\) \\
\hline 1644
1645 & \(\frac{\text { begin }}{\text { Validd }}\) dop \(=\) fals ; & 1745 & else error(19) \{ value not boolean \}; \\
\hline 1645
1646 &  & 1746 & 28, \({ }^{\text {if }}\) (ItsBoolean(value)) and (ItsBoolean(SFloa \\
\hline 1647 & then & 1747 & then \\
\hline 1648 & begin & 1748 & if (value \(=1.0\) ) or ( SFloat \(=1.0)\) \\
\hline 1649 & code := DOpTab[TokenTabPtr \({ }^{\text {² }}\). DOpIndx]. OpIndex; & 1749 & then SFloat \(:=1.0\) \\
\hline 1650 & hold : = TokenTabPtr; & 1750 & else SFloat \(:=0.0\) \\
\hline 1651 & TokenTabPtr : = TokenTabPtr^.NextToken; & 1752 & else error(19) \{ value not boolean \}; \\
\hline 1652 & if (code > 80) then ValidD := true & 1752 & 29: \\
\hline 1653 & else & 1754 & if value > SFloat \{ maximum or ceiling \} \\
\hline 1654 & if TokenTabPtr \({ }^{\text {- }}\).noun \(=\) SpecOper & 1755 & then SFloat := value; \\
\hline 1655 & then & 1756 & 30: \\
\hline 1656 & if SpecSymbol(XPeriod) & 1755 & if value < SFloat \{ minimum or floor \} \\
\hline 1657 & then & 1757 & then SFloat := value; \\
\hline 1658 & begin & 1758 & 31: \\
\hline 1659 & if TokenTabPtr^.noun \(=\) DyadOper & 1759 & if (value * SFloat) < 0.0 \\
\hline 1660 & then & 1761 & then error(50) \{ number and base of different sign \} \\
\hline 1661 & begin & 1762 & else \({ }^{\text {els }}\) ( \({ }^{\text {a }}\) \\
\hline 1662 & if DOpTab[TokenTabPtr \({ }^{\wedge}\). DOp Indx]. Op Index \(<=80\) & 1783 & \(\frac{\text { SFloat }}{}:=(\ln (a b s(S F l o a t))) /(\ln (\) abs \((\) value \())\) ) \\
\hline 1663 & then & 1764 & \{ log to a base \} \\
\hline 1664 & begin code : \(=\) code + (100 * DOpTab[TokenTabPtr \({ }^{\text {² }}\). & 1765 & end \{ case \} \\
\hline 1665 & code : \(=\) code + (100 * DOpTab[TokenTabPtr \({ }^{\text {² }}\). & 1766 & end \{ dyadcomp \}; \\
\hline 1666 & DOp Indx]. OpIndex); & 1767 & \\
\hline 1667 & hold : = TokenTabPtr; & 1768 & \\
\hline 1668 & TokenTabPtr \(:=\) TokenTabPtr \({ }^{\text {² }}\). NextToken; & 1769 & procedure IndexGenerator(arg: TypeValtabPtr); \\
\hline 1669 & Validd : \(=\) true & 1770 & \{ monadic iota operator \} \\
\hline 1670 & el end \({ }^{\text {ce }}\) error(27) , invalid inner product exp \} & 1771 & \\
\hline 1671 & else error(27) \{ 'invalid inner product exp \} & 1772 & var \\
\hline 1672 & el \(\frac{\text { end }}{\text { ce }}\) & 1773 & iotaIndex, TopValue: integer; \\
\hline 1673
1674 & \(\frac{\text { else }}{\text { if }}\) Tokentabptr \({ }^{\wedge}\). \({ }^{\text {noun }}=\) SpecOper & 1774 & \\
\hline 1674 & if TokentabPtrn.noun \(=\) Specoper & 1775 & begin \\
\hline 1675 & then \({ }^{\text {begin }}\) & 1776 & if arg^.dimensions <> 0 \\
\hline 1676
1677 & \(\frac{\text { begin }}{\text { if }}\) SpecSymbol (XLittlecircle) & 1777 & then error (21) \{ argument not a scalar \} \\
\hline 1677 & if Specsymbol(xLittlecircle) & 1778 & else \\
\hline 1678 & & 1779 & \(\frac{\text { if }}{\text { arg }}\). FirstValue^.RealVal \(<0.0\) \\
\hline 1679 & \(\frac{\text { begin }}{\text { end }}\) code \(:=10 \star\) code; Valid : \(=\) true & 1780 & then error(22) \{ argument is negative \\
\hline 1680 &  & 1781 & else \\
\hline 1681
1682 & else error(26) \({ }^{\text {en }}\) 'inval outer prod exp & 1782 & if (arg^.FirstValue^.RealVal) - (1.0 * trunc(arg`. \\
\hline 1682
1683 & el \(\frac{\text { end }}{\text { se }}\) error (26) \{ same as above \} & 1783 & FirstValue \({ }^{\text {. ReatVa()) }}\) <> 0.0 \\
\hline 1683
1684 & end else error(26) same as above & 1784 & then error(23) \{ argument is not an integer \} \\
\hline 1684
1685 & elsevalidd := true & 1785 & else \\
\hline 1685
1680 & else \({ }^{\text {else }} \mathrm{VatidD}:=\) true; & 1786 & begin \\
\hline 1687
1688 & \begin{tabular}{l}
end; \\
op := ValidD;
\end{tabular} & 1787
1788 & new(NewValtablink);
OldVal TabLink \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline 1789 & NewValTabLink \({ }^{\text {- }}\). NextValTabLink : \(=\) nil; & 1891 & procedure reduction(arg: TypeValTabPtr); \\
\hline 1790 & NewVal TabLink \({ }^{\text {. }}\). ForwardOrder \(:=\) true; & 1892 & procedure reduction(arg. TypeValtabptr) \\
\hline 1791 & NewValTabLink \({ }^{\text {² }}\). IntermedResult : \(=\) true; & 1893 & var \\
\hline 1792 & NewValTabLink^.dimensions : \(=1\) \{ result is a vector \}; & 1894 & counter, RowLength: integer; \\
\hline 1793 & new(NewDim); NewValTabLink \({ }^{\text {- }}\).FirstDimen \(:=\) NewDim; & 1895 & Sfloat: real; \\
\hline 1794 & TopValue : = trunc(arg^.FirstValue^. RealVal) & 1895 & \\
\hline 1795 & \{ last index generd \}; & 1897 & begin \\
\hline 1796 & NewDim \({ }^{\text {. }}\) dimenlength : \(=\) TopValue; & 1898 & if (arg^. \({ }^{\text {dimensions }}=0\) ) or ( arg \(^{\wedge}\). FirstValue \(=\) nil \()\) \\
\hline 1797 & NewDim^.NextDimen : = nil; iotaIndex := 1; & 1899 & then \\
\hline 1798 & switch := true; & 1900 & error(24) \{ argument is a scalar or vector of length zero \\
\hline 1799 & while iotaIndex <= TopValue do & 1901 & else \\
\hline 1800 & beg in & 1902 & if (arg^. \({ }^{\text {dimensions }}=1\) ) and (arg^.firstDimen \({ }^{\wedge}\). \({ }^{\text {dimentength }}\) \\
\hline 1801 & new(NewValues); NewValues^. Realval := iotaIndex; & 1903 & = 1) - \\
\hline 1802 & if switch & 1904 & then error(51) \{ argument is a vector of length one \} \\
\hline 1803 & then & 1905 & else \\
\hline 1804 & begin & 1906 & begin \\
\hline 1805 & switch := false; & 1907 & new(NewVa(TabLink) ; \\
\hline 1806 & NewValTablink \({ }^{\text {. }}\).FirstValue \(:=\) NewValues & 1908 & OldValTabLink \({ }^{\text {. Next ValTabLink }:=\text { NewValTabLink; }}\) \\
\hline 1807 & end & 1909 & NewVal TabLink \({ }^{\text {. }}\). NextValTabLink \(:=\) nil; \\
\hline 1808 & else NewVal Ptr \({ }^{\text {n }}\). NextValue \(:=\) NewValues; & 1910 & NewValTabLink \({ }^{\text {. }}\). IntermedResult \(:=\overline{\text { true }}\) \\
\hline 1809 & NewValPtr : = NewValues; & 1911 & if argn. ForwardOrder then ReverseLinkList(arg); \\
\hline 1810 & iotaIndex : \(=\) iotaIndex + 1 & 1912 & NewVal Tablink \({ }^{\text {n }}\).ForwardOrder \(:=\) false; \\
\hline 1811 & end; & 1913 & NewValTabLinkn.dimensions : \(=\arg ^{\wedge}\).dimensions - 1; \\
\hline 1812 & if switch & 1914 & DimPtr := arg^.FirstDimen; switch := true; \\
\hline 1813 & then & 1915 & white DimPtr \({ }^{-}\).NextDimen <> nil do \\
\hline 1814 & NewVal TabLink^.FirstValue : = nil & 1916 & begin \{ build dimensions of result \} \\
\hline 1815 & \(\{\) result is vector of length 0 \} & 1917 & new(NewDim) ; \\
\hline 1816 & else NewValues \({ }^{\text {² }}\). NextValue : \(=\underline{\text { nil }}\) & 1918 & if switch \\
\hline 1817 & end & 1919 & then \\
\hline 1818 & end \{ indexgenerator \}; & 1920 & begin \\
\hline 1819 & & 1921 & switch := false; \\
\hline 1820 & & 1922 & NewValTabLink \({ }^{\text {® }}\).FirstDimen : \(=\) NewDi \\
\hline 1821 & procedure ravel(arg: TypeValTabPtr); & 1923 & end \\
\hline 1822 & \{ monadic comma operator \} & 1924 & else NewPtr \({ }^{\text {a }}\). NextDimen \(:=\) NewDim; \\
\hline 1823 & & 1925 & NewDim.dimenlength : \(=\) DimPtr\({ }^{\text {a }}\).dimenlength; \\
\hline 1824 & var & 1926 & NewPtr := NewDim; DimPtr := DimPtr^.NextDimen \\
\hline 1825 & elements: integer; & 1927 & end; \\
\hline 1826 & & 1928 & if switch \\
\hline 1827
1828 & \begin{tabular}{l}
begin \\
new(NewValTabLink);
\end{tabular} & 1929 & \(\frac{\text { then }}{\text { NewValTabLinkn }}\). FirstDimen \(:=\) nil \\
\hline \[
\begin{aligned}
& 1828 \\
& 1829
\end{aligned}
\] & new(NewValtabLink); & 1930 & NewValTabLink^. FirstDimen : \(=\) nil \\
\hline 1830 & NewValTabLink \({ }^{\text {- }}\). NextValTabLink : \(=\) nil; & 1932 & else NewDim \({ }^{\text {. }}\) NextDimen \(:=\) nil; \\
\hline 1831 & NewValTabLink \({ }^{\text {- }}\). IntermedResult \(:=\) true; & 1933 & RowLength := Dimptri.dimenlength; \\
\hline 1832 & NewValTabLink^. ForwardOrder : = argn..Forwardorder; & 1934 & ValPtr : = arg \({ }^{\text {. }}\) FirstValue; switch : \(=\) true; \\
\hline 1833 & NewValTabLink \({ }^{\text {- }}\) dimensions \(:=1\) \{ result is a vector \}; & 1935 & while ValPtr <> nil do \\
\hline 1834 & new(NewDim); NewValTabLink^.FirstDimen : NewDim; & 1936 & begin \{ perform reduction\} \\
\hline 1835 & NewDim^.NextDimen : \(=\) nil; switch := true; & 1937 & SFloat : \(=\) ValPtr \({ }^{\text {n }}\). RealVal \\
\hline 1836 & ValPtr : \(=\) arg \({ }^{\text {- }}\) FirstValue; elements \(:=0\); & 1938 & \{ sfloat gets last value in row \}; \\
\hline 1837 & while ValPtr <> nil do & 1939 & ValPtr \(:=\) ValPtr \({ }^{\text {® }}\). NextValue; \\
\hline 1838 &  & 1940 & for counter := 2 to RowLength do \\
\hline 1839 & new(NewValues); NewValues \({ }^{\text {a }}\). RealVal : \(=\) ValPtr \({ }^{\text {n }}\).RealVal; & 41 & begin Dy \({ }^{\text {Data }}\) ( \\
\hline 1840 & elements := elements + 1; & 1942 &  \\
\hline 1841 & if switch & 1943 & ValPtr : = ValPtr \({ }^{\text {- }}\).NextValue \\
\hline 1842 & then & 1944 & \(\frac{\text { end }}{\text { nem }}\) (NewValues) . NewValues^. Realval : \(=\) Sploat, \\
\hline 1843 & begin & 1945 & new(NewValues); NewValues^. Realval := SFloat; \\
\hline 1844 & switch := false; & 1946 & if switch \\
\hline 1845 & NewValTabLink^.FirstVatue := NewValues & 1947 & \(\frac{\text { then }}{\text { begin }}\) \\
\hline 1846 & elend \({ }^{\text {end }}\) NewValPtr^ . NextValue : \(=\) NewValues; & 1948 & \(\frac{\text { begin }}{\text { switch }}:=\) false; \\
\hline 1847
1848 &  & 1949
1950 & Switch : \(=\) false;
NewValTabLink".FirstValue : \\
\hline 1849 & end; & 1951 & end \\
\hline 1850 &  & 1952 & else NewValPtr \({ }^{-}\). NextValue : \(=\)NewValues; \\
\hline 1851 & if switch then NewValTabLink \({ }^{\text {n }}\).FirstValue \(:=\underline{\text { nil }}\) & 1953 & NewValPtr : = NewValues \\
\hline 1852 & endse NewValues \({ }^{\text {a }}\). NextValue : \(=\) nil & 1954 & end; \\
\hline 1853 & end \{ ravel \}; & 1955 & NewValues^. NextValue := nil \\
\hline 1854 & & 1956 & \\
\hline 1855 & & 1957 & end \{ reduction \}; \\
\hline 1856 & procedure Shape Of(arg: TypeVa(TabPtr); & 1958 & \\
\hline 1857 & \{ monadic rho operator \} & 1959 & procedure monadic(arg: TypeValTabPtr; token: TokenPtr); \\
\hline 1859 & begin & 1960 & ( operations with codes between 1 and 31 \} \\
\hline 1860 & new(NewVa(TabLink); & 1962. & \\
\hline 1861 & OldValTabLink \({ }^{\text {. }}\). NextValTabLink \(:=\) NewValTabLink; & 1963 & begin \\
\hline 1862 & NewValTabLink^. NextValTabLink := nil; & 1964 & if token^.noun \(=\) ReductOper then reduction(arg) \\
\hline 1863
1864 & \begin{tabular}{l}
NewValTabLink^. IntermedResult := true; \\
NewValTabLink^.ForwardOrder := true;
\end{tabular} & 1965 & if code > 20 \\
\hline 1864
1865 &  & 1966 & \[
\begin{aligned}
& \text { if code > } 20 \\
& \text { then }
\end{aligned}
\] \\
\hline 1865 & new(NewDim); NewDim \({ }^{\text {a }}\). dimenlength \(:=\mathrm{arg}^{\text {a }}\).dimensions; & \[
\begin{aligned}
& 1967 \\
& 1968
\end{aligned}
\] & case code of \\
\hline 1867 & NewValTabLink \({ }^{\wedge}\).FirstDimen \(:=\) NewDim; & 1969 & 21: IndexGenerator(arg); \\
\hline 1868 & NewDim^NextDimen := nil; switch := true; & 1970 & 22: Shape Of (arg); \\
\hline 1869 & DimPtr \({ }^{\text {a }}=\mathrm{arg}^{\text {a }}\).FirstDimen; & 1971 & 23: ravel(arg) \\
\hline 1870 & while Dimptr <> nil do & 1972 & \(\frac{\text { end }}{\text { Le }}\) \{ case \} \\
\hline 1871 & begin \{ argument dimensions become result values \} & 1973 & beg in \\
\hline 1872 & new (NewVa Lues) ; & 1974 & new(NewVa(TabLink); \\
\hline 1873 & NewValues \({ }^{\text {n }}\).RealVal : \(=\) DimPtr \({ }^{\text {n }}\).dimenlength; & 1975 & OldValTabLink \({ }^{\text {- }}\) NextValTabLink : \(=\) NewValTabLink; \\
\hline 1874 & if switch & 1976 & NewValTabLink \({ }^{\text {. }}\) NextValTabLink \(:=\) nil; \\
\hline 1875
1876 & then & 1977 & NewValTablink \({ }^{\text {. }}\) IntermedResult : \(=\) true; \\
\hline 1877 & \(\frac{\text { begin }}{\text { switch }}\) : \(=\) false; & 1979 & NewValTabLink^. ForwardOrder := arg^. ForwardOrder; \\
\hline 1878 & NewValTabLink \({ }^{\text {² }}\). \({ }^{\text {irstValue }}\) : \(=\) NewValues & 1980 &  \\
\hline 1879 & end & 1981 & while DimPtr <> nil do \\
\hline 1880 & else NewValPtr \({ }^{\text {a }}\). NextValue \(:=\) NewValues; \({ }^{\text {den }}\), & 1982 & \(\frac{\text { begin }}{}\) \{ duplicate \(\frac{1}{\text { dimensions of arg into result \} }}\) \\
\hline 1881 & NewValPtr : \(=\) NewVatues; DimPtr : \(=\) DimPtr^ .NextDimen & 1983 & new(NewDim); \\
\hline 1882 & if end; & 1984 &  \\
\hline 1883 & if switch & 1985 & if switch \\
\hline 1884 & \(\frac{\text { then }}{\text { NewValTabLink }}\). FirstValue \(:=\) nil & 1987 & then \\
\hline 1885
1886 & NewValTabLink \({ }^{\text {a }}\). FirstValue \(:=\frac{\text { nil }}{\text { a }}\) ( result is a vector of length 0\(\}\) & 1988 & \(\frac{\text { begin }}{\text { switch }}\) : \(=\) false: \\
\hline 1887 & else NewValues \({ }^{\text {. }}\). \({ }^{\text {axtValue }}:=\underline{\text { nil }}\) & 1989 & NewValTabLink \({ }^{\text {c }}\). \(\mathrm{FirsstDimen}^{\text {a }}=\) NeWDim \\
\hline 1888 & end t shapeof 1; & 1990 & end \\
\hline
\end{tabular}

\begin{tabular}{|c|c|}
\hline 2093 & begin \\
\hline 2094 & switch := false; \\
\hline 2095 & NewValtablink \({ }^{\text {- }}\) Firstvalue : \(=\) NewValues \\
\hline 2096 & end \\
\hline 2097 & else NewValPtr \({ }^{\text {- }}\). NextValue : \(=\) NewValues; \\
\hline 2098 & \(\overline{\text { NewValues }}\). Realval : \(=\) LeftValPt \(r^{\wedge}\). RealVal; \\
\hline 2099 & NewValPtr : \(=\) NewValues; \\
\hline 2100 & LeftValPtr : = LeftValPtr \({ }^{\text {n }}\).NextValue \\
\hline 2101 & end; \\
\hline 2102 & RightValPtr : = Rightarg^.firstValue; \\
\hline 2103 & while RightvalPtr <> nil do \\
\hline 2104 & begin ( transfer right arg values (if any) \\
\hline 2105 & new(NewValues); \\
\hline 2106 & if switch \\
\hline 2107 & then \\
\hline 2108 & begin \\
\hline 2109 & switch := false; \\
\hline 2110 & NewValtabLink \({ }^{\text {. FirstValue }}\) : \(=\) NewValues \\
\hline 2111 & end \\
\hline 2112 & else NewValPtr \({ }^{-}\). NextValue : \(=\)NewValues; \\
\hline 2113 & \(\overline{\text { NewValues }}\). RealVal : \(=\) RightValPtr \({ }^{\wedge}\). RealVal; \\
\hline 2114 & NewValPtr : \(=\) NewValues; \\
\hline 2115 & RightValPtr : \(=\) RightValPtr \({ }^{\text { }}\). NextValue \\
\hline 2116 & end; \\
\hline 2117 & NewValues^.NextValue : = nil \\
\hline 2118 & end \{ transfer of values \} \\
\hline 2119 & end \\
\hline 2120 & end ( catenate ); \\
\hline 2121 & \\
\hline 2122 & \\
\hline 2123 & procedure IndexOf(LeftArg, RightArg: TypeValtabPtr); \\
\hline 2124 & \{ dyadic iota operator \} \\
\hline 2125 & var \\
\hline 2127 & MapIndex, icount, Testlength, OneMore: integer; \\
\hline 2128 & \\
\hline 2129 & begin 1 index of \\
\hline 2130 & if LeftArg*.dimensions <> 1 \\
\hline 2131 & then error(29) \{ left argument is not a vector \\
\hline 2132 & else \\
\hline 2133 & begin \\
\hline 2134 & new(NewVa(TabLink); \\
\hline 2135 & OLdValTabLink \({ }^{\text {a }}\). NextValTabLink \(:=\) NewValTabLink; \\
\hline 2136 & NewVal Tablink . NextValTabLink := nil; \\
\hline 2137 & NewValTabLink \({ }^{\text {. }}\) IntermedResult \(:=\) true; \\
\hline 2138 & if not LeftArg*.Forwardorder \\
\hline 2139 & then ReverseLinkList(LeftArg); \\
\hline 2140 & NewValTabLink^. ForwardOrder : = RightArg^. Forwardorder; \\
\hline 2141 & NewValTabLink \({ }^{\wedge}\). dimensions : \(=\) RightArg \({ }^{\text {² }}\).dimensions; \\
\hline 2142 & if RightArg*.dimensions \(=0\) \\
\hline 2143 & then \\
\hline 2144 & NewValTabLink \({ }^{\text {a }}\).FirstDimen \(:=\) nil \\
\hline 2145 & \{ right argument is a scalar\} \\
\hline 2146 & else \\
\hline 2147 & begin \(\{\) build dimensions of result \} \\
\hline 2148 & switch := true; Dimptr := RightArg^.FirstDimen; \\
\hline 2149 & while DimPtr <> nil do \\
\hline 2150 & begin \\
\hline 2151 & new(NewDim); \\
\hline 2152 & if switch \\
\hline 2153 & \(\frac{\text { then }}{\text { begin }}\) \\
\hline 2154 & \(\frac{\text { beg in }}{\text { switch }}:=\) false; \\
\hline 2155
2156 & switch \(:=\) false;
NewValTabLink \\
\hline 2157 & end \\
\hline 2158 & else NewPtr \({ }^{-}\). NextDimen : \(=\)NewDim; \\
\hline 2159 & NewDim^.dimenlength : \(=\) DimPtr \({ }^{\text {n }}\).dimenlength; \\
\hline 2160 & NewPtr : \(=\) NewDim; DimPtr \(:=\) DimPtr \({ }^{\text {- }}\).NextDimen \\
\hline 2161 & end; \\
\hline 2162 & NewDim \({ }^{\text {a }}\).NextDimen \(:=\) nil \\
\hline 2163 & switch \(\frac{\text { end }}{}=\) true; RightValPtr \(:=\) RightArgn.FirstValue; \\
\hline 2164 & switch := true; RightValPtr := RightArg. FirstValue; white RightValPtr <> nil do \\
\hline 2165
2166 & begin \\
\hline 2167 & new(NewValues); \\
\hline 2168 & if switch \\
\hline 2169 & then \\
\hline 2170 & \(\frac{\text { begin }}{\text { switch }}\) : \(=\) false. \\
\hline 2171 & Switch := false; \\
\hline 2172 & NewValTablink . FirstValue : \(=\) NewValues \\
\hline 2173 &  \\
\hline 2174 & icount \(:=1\); LeftValPtr \(:=\) LeftArgi.FirstValue; \\
\hline 2175
2176 & TestLength : = Leftarg^.FirstDimen \({ }^{\text {- }}\). \({ }^{\text {dimenlength }}\) \\
\hline 2177 & \{length of left arg \}; \\
\hline 2178 & OneMore \(:=\) TestLength +1 \\
\hline 2179 & MapIndex \(:=\) OneMore; \({ }^{\text {a }}\) len plus one ]; \\
\hline 2180 & while (icount <= TestLength) and (MapIndex = OneMore) do \\
\hline 2181 & begin \\
\hline 2182 & try to match value in right arg with one in left arg \} \\
\hline 2184 & if LeftValPtr \({ }^{\text {n }}\).RealVal \(=\) RightValPtr \({ }^{\boldsymbol{n}}\). RealVal \\
\hline 2185 & then MapIndex \(:=\) icount \{ value match \}; \\
\hline 2186 & \begin{tabular}{l}
icount := icount + 1; \\
LeftValPtr : = LeftValPtr". NextValue
\end{tabular} \\
\hline 2187 & end; \\
\hline 2188 & NewValues^. Realval : = MapIndex; \\
\hline 2189 & NewValPtr : = NewValues; \\
\hline 2190 & RightValPtr : = RightValPtr \({ }^{\wedge}\). NextValue \\
\hline 2191 & end \\
\hline 2193
2194 & \{ if no match, index becomes one more than length of left arg \} ; \\
\hline
\end{tabular}

new(NewVa(TabLink);
OldValTabLink^. NextValTabLink := NewValTabLink;
NewValTabLink^. NextValTabLink := nil;
NewValTabLink^. IntermedResult := \(\begin{aligned} & \text { true; }\end{aligned}\)
if not LeftArg \({ }^{\text {n }}\). ForwardOrder
then ReverseLinkList (LeftArg);
if not RightArgn. ForwardOrder
then ReverseLinkList (RightArg) ;
NewValTabLink^. ForwardOrder := true;
NewValTabLink^.dimensions := LeftArg^. .dimensions + RightArg^.dimensions - 2;
if NewValTabLink^. dimensions < 0
then NewValTabLink^.dimensions \(:=0\);
switch := true; LastLeftDim:=0;
if LeftArg^.FirstDimen <> nil
then
begin \{ copy all but last of left arg dims into result \}
LeftSkip \(:=1\); DimPtr \(:=\) LeftArg \({ }^{\text {. FirstDimen; }}\)
while DimPtrn. NextDimen <> nil do
\(\frac{\text { begin }}{\text { new (NewDim) ; }}\) ( copt arg dimensions new(NewDim),
NewDim^. dimenlength := DimPtr^.dimenlength; LeftSkip:=LeftSkip * DimPtr^.dimenlength; if switch
\(\qquad\)
beg in
switch := false;
NewValTabLink^. FirstDimen := NewDim end
else NewPtr". NextDimen := NewDim; NewPtr : = NewDim; DimPtr := DimPtr^. NextDimen end;
LastLeftDim := DimPtr \({ }^{\wedge}\).dimenlength
end;
if RightArg^. FirstDimen <> nil
then
beg in
\{copy all but first of right arg dims into result \}
RightSkip: \(=1\);
DimPtr := RightArg^.FirstDimen^. NextDimen;
while Dimptr <> nil do
begin \{ copy right \(\frac{\text { arg dimensions \}}}{}\)
new(NewDim);
NewDim^. dimenlength : = DimPtr^. .dimenlength;
RightSkip := RightSkip * DimPtr^.dimenlength;
if switch
if sw
\(\qquad\)
switch : = false;
NewValTabLink^. FirstDimen := NewDim
end
el \(\overline{\text { se }}\) NewPtr \(r^{\wedge}\). NextDimen : \(=\) NewDim;
NewPtr : = NewDim; DimPtr := DimPtr".NextDimen end
end;
if switch then NewValTabLink^. FirstDimen : = nil
else NewDim - . NextDimen := nil;
\(\overline{\overline{i f} \text { LeftArg^. FirstValue }=n \overline{i l}}\) then LeftSkip \(:=0\);
if RightArg. FirstValue \(=\frac{\text { nil }}{\text { nil }}\) then RightSkip \(:=0\);
switch := true;
if RightArg^. FirstDimen <> nil
then FirstRightDim :=Right \(\overline{A r g}\). .FirstDimen \({ }^{-}\).dimenlength
eไse FirstRightDim :=0;
if FirstRightDim > LastLeftDim
then CommonLength \(:=\) FirstRightDim
else CommonLength \(:=\) LastLeftDim;
icount :=0; LeftValPtr := LeftArg*. FirstValue; while icount < LeftSkip do
begin \{loop for each row in left arg \}
lptr := LeftValPtr \{ hold start of row position \};
jcount :=0;
while jcount < RightSkip do
begin \{ loop for each column in right arg \}
LeftValPtr : = lptr;
RightValPtr : = RightArg^. FirstValue;
Icount := 0;
while lcount < jcount do
begin \(\{\) skip to starting value in right arg \}
RightValPtr := RightValPtr". NextValue;
if RightValPtr \(=\) nil then
RightValPtr : = \(\overline{\text { RightArg }}\). FirstValue
\{ extend arg \};
L count \(:=\) lcount +1
end;
count := 0;
while kcount < Commonlength do
\(\frac{\text { begin }}{\text { SFl }}\) loop for each element in row'column \}
SFloat : = RightValPtr \({ }^{\wedge}\). RealVal;
DyadComp(SFloat, LeftValPtrn.RealVal,
Inpro2Code);
value := SFloat
if kcount \(=0\)
then
set identity value for first time through \}
case Inpro1Code of
52, 53, 78: SFloat \(:=0.0\);
54, 55, 56, 77: SFloat \(:=1.0\);
71, 72, 73, 74, 75, 76: \{ null case \}
\(\qquad\)
lse Sfloat \(:=\) hold
DyadComp(SFloat, value, Inpro1Code); hold :=SFloat \{ save summer result \};
\begin{tabular}{|c|c|c|c|}
\hline 2399 & LeftValPtr : = LeftValPtr \({ }^{\text {- }}\).NextValue; & 2501 & procedure dyadic(LeftArg, RightArg: TypeValtabPtr); \\
\hline 2400 & if LeftValPtr \(=\) nil \({ }^{\text {then }}\) & 2502 & T operators with codes of 52 and higher \} \\
\hline 2401 & LeftValPtr : = Left Arg \(^{*}\).FirstValue & 2503 & \\
\hline 2402 & \{ extend arg \}; & 2504 & var \\
\hline 2403 & mcount : = 0; & 2505 & compatible: Boolean; \\
\hline 2404 & while mcount < RightSkip do & 2506 & arg: TypeValTabPtr; \\
\hline 2405 & begin \{ skip to next value in right arg \} & 2507 & Sfloat: real; \\
\hline 2406 & mcount := mcount + 1; & 2508 & \\
\hline 2407 & RightValPtr : = RightValPtr^. NextValue; & 2509 & begin \\
\hline 2408 & if RightValPtr \(=\) nil & 2510 & if code > 1000 then InnerProduct(LeftArg, RightArg) \\
\hline 2409 & then RightValPtr \(:=\) RightArg".FirstValue; & 2511 & else \\
\hline 2410 & end; & 2512 & \(\frac{\text { if }}{}\) code > 100 then OuterProduce(LeftArg, Rightarg) \\
\hline 2411 & \(k\) count : \(=\) kcount + & 2513 & else \\
\hline 2412 &  & 2514 & if code > 80 \\
\hline 2413 & new(NewValues); NewValues*.RealVal := Sfloat; & 2515 & then \\
\hline 2414 & if switch & 2516 & case code of \\
\hline 2415 & then & 2517 & 87: IndexOf(LeftArg, RightArg); \\
\hline 2416 & begin & 2518 & 88: reshape(LeftArg, RightArg); \\
\hline 2417 & switch : \(=\) false; & 2519 & 89: catenate(LeftArg, RightArg) \\
\hline 2418 & NewValTabLink^.FirstValue := NewValues & 2520 & end \{ case \} \\
\hline 2419 & end & 2521 & else \\
\hline 2420 & else NewValPtr^.NextValue : \(=\) NewValues; & 2522 & begin \(\{\) simple dyadics \\
\hline 2421 & NewValPtr : = NewValues; jcount := jcount + 1; & 2523 & compatible := true; \\
\hline 2422 & end; & 2524 &  \\
\hline 2423 & icour't := icount + 1 & 2525 & dimensions \(>=1\) ) \\
\hline 2424 & end; & 2526 & then \({ }_{\text {if }}\) \\
\hline 2425 &  & 2527 & if LeftArg^.dimensions <> RightArg^.dimensions \\
\hline 2426 & else NewValues \({ }^{\text {² }}\). NextValue \(:=\) nil & 2528 & \\
\hline 2427 & end & 2529 & \begin{tabular}{l}
compatible := false \\
\{ different ranks'neither scalar \}
\end{tabular} \\
\hline 2428
2429 & end \{ innerproduct \}; & 2530 & else \{ different ranks'neither scalar \} \\
\hline 2429
2430 & & 2531 & \(\frac{\mathrm{e} \text { (se }}{\text { begin }}\) \{ ranks match - check lengths ) \\
\hline 2431 & procedure OuterProduce(LeftArg, Rightarg: TypeVa(TabPtr); & 2533 & LeftDimptr : = LeftArg*.FirstDimen; \\
\hline 2432 & & 2534 & RigthDimPtr \(:=\) RightArg*.FirstDimen; \\
\hline 2433 & var & 2535 & while LeftDimptr <> nil do \\
\hline 2434 & OutProCode: integer; & 2536 & beg in \\
\hline 2435 & Sfloat: real; & 2537 & if LeftDimptra.dimenlength <> RigthDimPtr \({ }^{\text {a }}\). \\
\hline 2436 & & 2538 & dimenlength \\
\hline 2437 & beg in & 2539 & then \\
\hline 2438 & OutProCode : = code div 10; new(NewValTabLink); & 2540 & compatible : \(=\) false \{ different length(s) \}; \\
\hline 2439 & OldVal TabLink^. NextVal TabLink := NewValTabLink; & 2541 & LeftDimPtr : = LeftDimPtr \({ }^{\text {a }}\).NextDimen; \\
\hline 2440 & NewValTabLink . NextValTabLink := nil; & 2542 & RigthDimPtr : \(=\) RigthDimPtr \({ }^{\text {a }}\).NextDimen \\
\hline 2441 & NewVal TabLink \({ }^{\wedge}\). IntermedResult : \(=\) true; & 2543 & end \\
\hline 2442 & if not LeftArg*. Forwardorder & 2544 & end; \\
\hline 2443 & then ReverseLinkList (LeftArg); & 2545 & if compatible \\
\hline 2444 & if not RightArg \({ }^{\text {a }}\). Forwardorder & 2546 & \{ arguments suitible for dyadic operation \\
\hline 2445 & then ReverseLinkList(RightArg); & 2547 & then \\
\hline 2446 & NewVal TabLink^. Forwardorder : = true; & 2548 & begin \{ build dimensions of result \} \\
\hline 2447 & NewVal TabLink^. \({ }^{\text {dimensions }}\) : \(=\) LeftArg^. \({ }^{\text {dimensions + Right Arg^. }}\) & 2549 & if Rightarg \({ }^{\text {a }}\). dimensions \(>\) LeftArg \({ }^{\text {. }}\). \({ }^{\text {dimensions }}\) \\
\hline 2448 & dimensions; & 2550 & \(\frac{\text { then }}{\text { else }}\) arg : \(=\) RightArg \\
\hline 2449 & switch : = true; Dimptr := LeftArg*.FirstDimen; & 2551 &  \\
\hline 2450 & while DimPtr <> nil \(\frac{\text { do }}{\text { a }}\) & 2552 & arg : = LeftArg \{ result has shape of larger arg \}; \\
\hline 2451 & begin \{copy left arg dimensions to result \} & 2553 & new(NewVa(TabLink); \\
\hline 2452 &  & 2554 & OldValTabLink^.NextValTabLink := NewValTabLink; NewValTabLink*.NextValTabLink := nil; \\
\hline 2453 & if switch & 2555 & \begin{tabular}{l}
NewValTabLink^. NextValTabLink := nil; \\
NewValTabLink^. IntermedResult := true;
\end{tabular} \\
\hline 2454 & \(\frac{\text { then }}{\text { beg in }}\) & 2556
2557 & \begin{tabular}{l}
NewValTablink^. IntermedResult := true; \\
if LeftArg^. ForwardOrder <> RightArg*. ForwardOrder
\end{tabular} \\
\hline 2455
2456 &  & 2558 & then ReverseLinkList(LeftArg); \\
\hline 2457 & end & 2559 & NewValTabLink \({ }^{\text {n }}\). ForwardOrder \(:=\arg ^{\wedge}\). ForwardOrder; \\
\hline 2458 & else NewPtr \({ }^{\text {- }}\). NextDimen \(:=\) NewDim; & 2560 & NewValTabLink \({ }^{\wedge}\).dimensions \(:=\) arg^ \({ }^{\text {. }}\). dimensions; \(^{\text {a }}\) \\
\hline 2459 & NewPtr := NewDim; DimPtr \(:=\) DimPtr \({ }^{\text {a }}\).NextDimen & 2561 & switch := true; DimPtr := arg*.FirstDimen; \\
\hline 2460 &  & 2562 & while DimPtr <> nil \(\frac{\text { do }}{\text { a }}\), \\
\hline 2461 & Dimptr : = RightArg*.FirstDimen; & 2563 & begin (copy dimensions to result \} \\
\hline 2462 & While DimPtr <> nil do & 2564 & new(NewDim); \\
\hline 2463 & begin \{ copy dimensions of right arg to result \} & 2565 & NewDim^. dimenlength : \(=\) DimPtr \({ }^{\text {a }}\).dimenlength; \\
\hline 2464 & new(NewDim); NewDim \({ }^{\text {a }}\).dimenlength : \(=\) DimPt \({ }^{\text {a }}\). dimenlength;
if & 2566
2567 & if switch \\
\hline 2465 & \(\frac{\text { then }}{\text { thitch }}\) & 2567
2568 & \(\frac{\text { then }}{\text { beg in }}\) \\
\hline 2466 & & 2568
2569 & \\
\hline 2467
2468 & \(\frac{\text { begin }}{\text { switch }}:=\) false; NewValtablink . Firstoimen : NewDim & \(\begin{array}{r}2569 \\ \hline 250\end{array}\) & NewValTabLink^. FirstDimen := NewDim \\
\hline 2468
2469 & end Swtch := false; Newvaltablink .Firstoimen := NewDim & 2570 & end \\
\hline 2469 & else NewPtr \({ }^{\text {² }}\). NextDimen \(:=\) NewDim; & 2571 & else NewPtr \({ }^{\wedge}\). NextDimen : \(=\) NewDim; \\
\hline 2470
2471 & NewPtr : = NewDim; DimPtr : = DimPtr \({ }^{\text {- }}\).NextDimen & 2573 & \(\overline{\text { NewPt } r ~} \mathrm{t}\) : NewDim; \\
\hline 2471 & end; & 2574 & DimPtr : \(=\) DimPtr \({ }^{\text {a }}\).NextDimen \\
\hline 2473 & if switch then NewValTabLink \({ }^{\text {a }}\).FirstDimen : \(=\) nil & 2575 & if end; \\
\hline 2474 &  & 2576 & if switch \\
\hline 2475 & switch : = true; LeftValPtr : = Leftarg^.FirstValue; & 2577 & \(\frac{\text { then }}{\text { NewValTablink }}\) - FirstDimen : \(=\) nil \\
\hline 2476 & While LeftValPtr < \({ }_{\text {begin }}\) & 2578 & NewValTabLink . FirstDimen := nil \{result is a scal \} \\
\hline 2477 & \(\frac{\text { RightValPtr }}{}\) : \(=\) RightArg \({ }^{\text {a }}\).FirstValue; & 2579 & \\
\hline 2478 & RightvalPtr : \(=\) RightArg \({ }^{\text {a }}\).FirstValue;
while Right ValPtr <> nil do & 2580 & switch := true; \\
\hline 2479 & \(\frac{\text { while }}{\text { begin }}\) RightvalPtr <> nil do & 2581 & RightValPtr : \(=\) RightArg^.FirstValue; \\
\hline 2480 & \(\frac{\text { Sfloat }}{}\) := RightValPtr \({ }^{\text {a }}\).RealVal; & 2582
2583 & LeftValPtr : = Leftarg \({ }^{\text {. }}\) FirstValue; \\
\hline 2481
2482 & DyadComp(SFloat, LeftValPtr".RealVal, OutProCode); & \[
\begin{aligned}
& 2583 \\
& 2584
\end{aligned}
\] & ValPtr : = arg^.FirstValue; \\
\hline 2482 & new(NewValues); & 2584 & while Valptrer nit do \\
\hline 2483
2484 & if switch & 2586 & begin \{perform operation \} \\
\hline 2484 & then & 2587 & new( \(\mathrm{NewValues);}\) \\
\hline 2486 & beg in & 2588 & SFloat : = RightvalPtra.Realval; \\
\hline 2487 & Switch:= false; & 2589 & DyadComp(SFloat, LeftValPtr \({ }^{\text {a }}\). RealVal, code); \\
\hline 2488 & NewValTabLink^.FirstValue : \(=\) NewValues & 2590 & NewValues^.RealVal : = SFloat; \\
\hline 2489 & el \(\frac{\text { end }}{\text { se }}\) NewValPtr \({ }^{\text { }}\). NextValue \(:=\) NewValues; & 2591 & then \\
\hline 2490 & NewValues \({ }^{\text {N }}\).RealVal \(:=\) SFloat; NewValPtr \(:=\) NewValues; & 2592 & \(\frac{\text { beg in }}{}\) \\
\hline 2491 &  & 2593 & \(\frac{\text { segitch }}{\text { s }}=\) false; \\
\hline 2492 & & 2594 & NewValTabLink \({ }^{\text { }}\). FirstValue \(:=\) NewValues \\
\hline 2493 & LeftValPtr := LeftValPtr^. NextValue & 2595 & end \\
\hline 2494 & end; & 2596 & elsewValPtr \({ }^{\wedge}\). NextValue : \(=\) NewValues; \\
\hline 2495 & if switch then NewValTabLink^.FirstValue := nil & 2597 & NewValPtr := NewValues; \\
\hline 2496 & else NewValues^.NextValue : \(=\) nil & 2598 & ValPtr : = ValPtr \({ }^{\wedge}\). NextValue; \\
\hline 2497
2498 & end \(\{\) outerproduct \}; & 2599
2600 & LeftValPtr : = LeftValPtr \({ }^{\text {² }}\). NextValue; \\
\hline
\end{tabular}
2701
\(\frac{\text { begin }\{\text { parser }}{\text { assign }}==\)
\(\frac{\text { begin }\{\text { parser }\}}{\text { assign }:=\text { false; assign1 }:=~ f a l s e ; ~ D o n e P a r s e ~}:=\) false;
    repeat
\(\begin{array}{lll}2704 & \text { expression(Validexp) \{ checks for valid expression \}; } \\ 2705 & \text { if not Validexp then error(10) \{ 'invalid expression' \}}\end{array}\)
2704
2705
2706
2705
2706
2706
2707
        if not Validexp then error(10) \{ 'invalid expression' \}
            \(\frac{\text { if }}{\text { if }}\) SpecSymbol(XRight Arrow)
            if Sp
                if not ( (OperTabPtr \({ }^{\wedge}\). OperPtr \({ }^{\wedge}\).FirstValue \(=\) nil) and (
                \(\overline{\text { OperTabPtr }}\). OperPtr \({ }^{\wedge}\). dimensions \(>0\) ))
            then \{ branch \}
                \{result of expression is at opertabptr \}
                    if OperTabPtr \({ }^{\wedge}\). OperPtr \(r^{\wedge}\). FirstValue \({ }^{\wedge}\). RealVal - 1.0 * trunc
                    (OperTabPtr". OperPtrn. FirstValuen. RealVal) <> 0.0
                    then error(12) \{ stmt.num.to branch to not an integer \}
                    \(\frac{\text { then }}{\text { else }}\)
                    if SubrTabPtr \(=n i l\)
                    then
                    begin \{ function mode \}
                        TokenTabPtr := hold; DoneParse := true
                    end
                    \(\stackrel{\text { end }}{\text { end }}\)
                        trunc(OperTabPtr \({ }^{\wedge}\). OperPtrn. FirstValuen. RealVal) in
                    [1... (SubrTabPtr^. CalledSubr ". NumOf Statements)]
                    \(\frac{\text { then }}{\text { beg in }}\)
                    VFunchold := SubrTabPtr".CalledSubr^. FirstStatement;
                    for cnt : \(=1\) to trunc(OperTabPtr^. OperPtr \({ }^{\wedge}\).
                        FirstValue*.RealVal) do
                        begin
                            VFuncPtr : = VFunchold;
                                    TokenTabPtr :=VFuncPtr \({ }^{n}\). NextStmnt;
                                    VFunchold :=VFuncPtr \({ }^{-}\). NextVFunPrt
                                    end;
                    AuxOperTabPtr : = OperTabPtr;
                    OperTabPtr := OperTabPtr". LastOper;
                    dispose(AuxOperTabPtr); PtrLastOper := OperTabPtr;
                    TokenTabPtr : = VFuncPtr". NextStmnt
                    end
                    else ReturnToCallingSubr
            \(\frac{\mathrm{else}\{\text { successor \}}}{}\) \}
        else \{ successor \}
            beg in
                if not assign1 then OutPutVal; assign1 := false;
                if Subr TabPtr \(=n\) nit
                    then
                    begin \(\{\) interpretive \}
                    TokenTabPtr : = TokenTabPtr^. NextTokenं:
                    DoneParse := true
                \(\frac{\text { end }}{s e}\)
                else function \(\}\)
                    beg in
                    VFuncPtr : = VFuncPtra. NextVFunPrt;
                    DoneSuccessor := false;
                    repeat
                    if VFuncPtr <> nil
                    \(\frac{\text { then }}{\text { beg in }}\)
                        \(\frac{\text { begin }}{\text { TokenTabPtr }:=\text { VFuncPtr }}\). .NextStmnt;
                        DoneSuccessor := true
                    end
                    \(\frac{\text { beg in }}{\text { be }}\)
                            ReturntoCallingSubr:
                                    if TokenTabPtr \(r^{2}\), noun = StatEnd
                                    then DoneSuccessor \(:=\) true;
                                    endi \(\frac{\text { then }}{\text { end }}\)
                until DoneSuccessor;
                end;
                end
        until DoneParse;
        release \{ release memory \};
    end \{ parser \};
\(\left.\frac{\text { begin }}{\text { InitializeCharacterSet; ReadInErrorMsgs; }} \begin{array}{l}\text { scanner }\end{array}\right\}\)
    InitializeCharacterSet; ReadInErrorMsgs;
InitParser \{ initialize tables etc. \}; FillUpTables;
    InitParser \{ initialize tables etc. \}; FillUp
FunctionMode \(:=\) false; firstfunction \(:=\) true;
    OldValTablink \(:=\) nil; OldFuncTabPtr \(:=\) nil; OldVarTabPtr \(:=\) nil;
    OldTokenPtr \(:=\frac{n i l}{n i l}\) NewTokenPtr \(:=n i l ;\) NewFuncTabPtr \(:=n i l ;\)
    NewVFuncPtr \(:=\frac{n i l}{n i} ; \quad H o l d T o k e n P t r:=\frac{n i l}{n i} \quad\) TokenError \(:=\) false;
    NewValTablink \(:=\) nil; NewVarTabPtr \(:=\) nil; GetAPLstatement:
    while (APLstatement[1] <> character[ForwardSlash]) or (APLstatement[2]
        <> character[asterisk]) do \(f \quad\) * ends program
        begin
            \(\frac{\text { SkipSpaces; } \quad \text { TokenSwitch }:=\text { true; }}{}\)
            while (position \(<=\) LineLength) and (not TokenError) and (not
            LineTooLong) do
                    \(\frac{\text { begin }\{\text { scanning }\}}{\text { if } A P L s t a t e m e n t[p o s i t i o n] ~}=\) character[del]
                            f APLstatement[position] \(=\)
                    then \{ del encountered \}
                    if FunctionMode
                    if Fun
                    begin \{ end of current function \}
                                    \(\frac{\text { if }}{\text { New FuncTabPt } r ~<>~ n i l ~}\)
                                    Then NewFuncTabPtr \({ }^{\wedge}\). NumOf Statements \(:=\) FuncStatements;
                                    if. FuncStatements >0


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\section*{"Don't Fail Me Now"}

\author{
By Srully Blotnick
}

The government imposed a 55-mph speed limit on cars, not computers. Why, then, are computer owners going so slowly?

Are we in the early stages of a technology bust? Strange as this may sound at a time when the nation seems to have gone computer crazy, a good many scientists are starting to worry about just that.

Their concern stems from the massive switch in the computer business from a customer base consisting of a handful of large institutional buyers to millions of smaller ones. The computer finally has become a piece of mass-market electronics, much like video recorders. Why is that a problem? A basic rule of business is that risk accompanies opportunity. In this instance, the risk affects not only the companies in the field, but the entire country, thanks to the expanding economic importance of this industry. Its health will soon play a decisive role in determining the U.S.' international competitive position.

The risk in dealing with the mass market is always a simple one: The mob is fickle. What intrigues it today may leave it indifferent tomorrow. This time the fickleness could produce a national disaster. The U.S. has unwittingly invested a major portion of its capital - and even more important, its hopes - in this area. That's why some thoughtful workers in the field are beginning to pray quietly: "Don't fail me now."

How, specifically, do they see a failure occurring? The consensus view is as follows: "A Ferrari is exciting, but how exciting would it continue to be if the only place you could use it were your driveway? Well, that's exactly what is happening with way too many of the computers now being bought. Car or computer, people are eventually going to get tired of just looking at the thing and bragging about it to their friends. Then, the fad will pass. Computer manufacturing plants will close. Only a minuscule proportion of computer buyers are making good use of the machine's capabilities. They don't know enough about programming to make the machine really perform.',
'‘Well, suppose everyone learned BASIC?', I asked.

The overwhelming majority had a better idea: "BASIC is a very easy language to learn, but it would be enormously better, a dream come true, if everyone learned Pascal, which is far superior and just as easy to master.'

Dr. Srully Blotnick is a research psychologist and author of Getting Rich Your Own Way and Winning: The Psychology of Successful Investing.

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Since last summer I, therefore, have been collecting the opinions of everyone, from teachers and hobbyists to investors and small business owners, who know Pascal to see which books they consider best. A tally of the nearly 1,600 replies shows the following:

For people who know nothing at all about computers or computer programming, the best place to begin is R. Pattis' Karel the Robot: A Gentle Introduction to the Art of Programming (John Wiley, \$8.95). You don't need a computer to read this book (or the others about to be mentioned). By learning how to move a robot through the streets of a small town, you come to understand how programming instructs a computer to do what you want it to.

Pattis' book is about programming but doesn't actually teach the language. The elementary text that received the top rating in out survey was Arthur Keller's A First Course in Computer Programming with Pascal (McGraw-Hill, \$14.95). The book received high praise ('Very clear and easy to read'') from everyone from 17 to 70. It is suitable even as a high school text.

After Keller's book, the next step should be \(A\) Primer on Pascal by Conway, Gries and Zimmerman (Little, Brown, \$20). The consensus view: 'This book will help you deepen your understanding of the language once you've learned the elements." For those who already know BASIC, a good way to learn Pascal fast is Quick Pascal by D. Matuszek (John Wiley, \$11.95).

One work that was highly rated by advanced students was the second edition of Pascal - User's Manual and Report (Springer-Verlag, \$10.50) by K. Jensen and N. Wirth. That is hardly surprising since one of the coauthors, Nikolaus Wirth, invented the language.

To see what the language can really do, serious students will want to learn about data structures - that is, such things as lists, stacks, queues, trees, sets, records, recursion, sorting and searching. The three top-rated texts, all very well written, are: Data Structures and Algorithms by A. Aho, et al. (Addision-Wesley, \$28.95); Advanced Programming and Problem Solving with Pascal by G. Schneider and S. Bruell (John Wiley, \$26.95); and Data Structures Using Pascal by A. Tenenbaum and M. Augenstein (Prentis-Hall, \$25.95). As the authors of the first work comment, "The only prerequisite we assume is familiarity with some high-level programming language such as Pascal."

Finally, people with a background in probability theory rated the second edition of R. Cooper's Introduction to Queueing Theory (North-Holland Publishing Co., \$27) the best - clearest and most user-friendly book on the subject.

Summing up: Buying a computer and not learning to program it properly not only wastes money, it also stands a good chance of eventually harming the nation's economy.

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\title{
Computer Generated Population Pyramids Using Pascal
}

\author{
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}

\section*{Background}

During the past twenty years the development of computer applications in geography has been extensive. Hundreds of programs have been written and many are available to users through various dispensing institutions, particularly the Geography Program Exchange located at Michigan State University. \({ }^{1}\) Virtually all of the programs, however, are writen in FORTRAN and are suitable for easy installation primarily on large mainframe computers. \({ }^{2}\)

A similar situation exists in cartographic computer program development. Although the number and variety of programs written is extensive, the FORTRAN language is used almost exclusively, and the software is designed for use on large systems. A recent textbook in computer-assisted cartography provides only passing mention of microcomputer graphics in the field of cartography. \({ }^{3}\)

As a consequence of this situation, computer applications in geography and cartography are limited primarily to the larger colleges and universities that have mainframes and the faculty within the departments to teach the subjects. As a geographer in a small liberal arts college teaching not only introductory cartography but a course in micro-based computer mapping, I feel somewhat like a pioneer trying to make a clearing in the wood without the proper tools. The situation is further exacerbated because liberal arts colleges have not been as highly revered by the computer industry as have the high technology learning centers and are consequently not receiving anywhere near the number of equipment grants or the same degree of personnel support.

Yet, more than one writer has commented on the need for a closer association between the computer industry and the liberal arts college. In a recent editorial in Datamation, John L. Kirkely stated the following:

We urge our industry to work with the liberal arts colleges to develop courses of study that combine the humanities and the sciences. A merging of these artificially separated disciplines could be a powerful tonic for both our colleges and our corporations. \({ }^{4}\)

I believe that, in time, changes will be made which will result in the liberal arts colleges receiving their fair share of industry support. In the meantime, however, individuals in those colleges will continue to make contributions to the furtherance of computer applications in what would be considered today to be non-traditional disciplines. The set of programs included in this paper are suggestive of the kinds of things faculty can produce
and which 1) are effective vehicles for developing the understanding of key concepts in a discipline (demography in this case); 2) are produced with a cost factor reflecting only the programmer's time; 3) can be easily implemented on any system, micro to mainframe; and 4) are written in the programming language of the day, Pascal.

\section*{The Population Pyramid (Age Structure) Diagram}

It is abundantly clear that world population continues to grow at a less than acceptable rate, and that some regions, particularly those with countries exhibiting low levels of economic development, have exceptionally high rates of growth. \({ }^{5}\) The population pyramid is a useful diagram to study the composition of the population of any country or region.

In the diagram, age groupings of five years each ( \(0-4,5-9,10-14, \ldots\), up to \(75+\) ) are presented for both male and female segments of the population. The scale along the horizontal axis reports the percentage of the total population in each of the age groups. Generally a pyramid shape wide at the bottom (young age groups) is representative of a fast growing population while an age structure more evenly represented along the year's axis identifies a population that is stabilizing and that does not have a high rate of increase. The industrialized and urbanized countries in the developed world would fall into the later category; the less developed in the former.

\section*{The Programs}

Three programs have been developed for student use in an introductory human geography course. \({ }^{6}\) It is not necessary that the students know the Pascal language in order to run the programs. Introduction is given in class on login/logoff procedures and how to access the programs. The student need only find suitable information in the appropriate statistical source for each of the age groups for a particular region, round these values to a whole number, and enter the numbers in the sequence described in the program prompts.

The programs developed include:
1) pyramid_file - This program is used to create an external file of information including the region names, year of the data, and the percentages of male and female in each age group. Following the input, procedure echo-data publishes all the information entered for verification. If there were no input errors, the student selects the appropriate key and the program stores the information in an external file in the student's account. The listing of program pyramid_file follows:
program \(\operatorname{cyramidafile(input,~outout):~}\)

\section*{（a program to create an exterral the ot population pyramid tatal}
```

    const
    ```
        separator =
type
        data \(=\)
            cord
            country: packed arrav \(11 . . \mid\) bl char:
            year: packed array (1..4) of char:
            malenercent, femalepercent: array \((1 . .16)\) nt irteser
        end:
        identifiers \(=f 1 l e\) of तata;
var
        temp: data;
        infor identifiers:
        fllename: packed array (1..1n) of erar:
        answer: char:
procedure readadatal var．
：integer：
neain
＊riteln：
wite（＇enter flle name：＇）
eadin（filenare）：
writeln；
writec enter place name－－use it columns： ：）：
readin（temr．countrv）：
writelns
writeln：
write（ enter the year of the data：＊）：
readin（temp．year）：
writeln：
writeln（ now，enter rale and female rercertages for＇）：

writeln（＂frof the nighest ace aroup to tre lowest．＂）：
writeln（ you must nave 16 entries for eacn orour．＂）；
writelncenter as fnteners all on the sate line．＂；
writeln：
for \(1:=1\) to 10 do
read（temp．malenercent（1））；
writeln（ \({ }^{\prime}\) next，enter the female rercentaces．＂）；
for \(1 \quad i=1\) to 10 do
read（temr，femalerercent（i）） readin：
end \{reat-data\} :
procedure echo－data；
var
1：integer；
beain
writelng
writeln（separator）：
writeln；
writeln（．the followina information was enteredi＂）：
writeln；
writeln（, ，terr．country）
writeln；
writein（e．，temr．year）：
writeln：
writel rale
tor \(1:=1\)
tor \(1:=1\) to 16 do
write（terp．ralefercent（i）： 3\()\)
writeln；
writeln：
ritec female：＇）：
or \(1:=1\) to 16 to
write（temu．tenalefercent（il：3）：
writeln：
Mriteln（sedarator）：
writeln（ is the iptormation correct？＂））

writeln（．．．if yes，enter
readin（answer）：
writelnt
and（ecno－datal ：
procedure storeadata；
var
1：integer
reain
rewrite（info，fllename）；
infon．country iz terc，country：
rut（info）：
infon，year ：＝temf．year；
put（info）：
or \(1:=1\) to 1000
healn
infoc．
info \(0^{\circ}\) ralepercent（1）\(:=\) temp．nalepercent \((1)\)
rut（info）
end：
or \(1:=1\) to 10 तo
reain femalenercent（1）\(:=\) temo．femalefercent（1）：
put（info）
end：
close（1nto）
end \｛storeatata\} :
```

begin
readadata;
echo_data:
1f (answer = 'y') or (answer = 'y') then
ca!
store_dतta;
writeln(. information storecin the flle: !, tilename)
end
lse
nriteln(* invalid data: run the nroarar acalm")
nd.

```

2）get＿pyr＿file－An editing program which the student may use to access an external file，display the contents，and make any necessary changes．This pro－ gram would come in handy if more recent data is re－ ceived and the file is to be updated．The listing or program get＿pyr＿file and an example run of informa－ tion contained in the external file，SWEDEN．PYR， follow：
```

program aet_nyr-file(incut, outwut):
la proaram to examine the external data flle
created oy the program, 'oyramid_flle', and
to make changes if necessaryt
type
data =
record
country: packed array (1..15] nt char:
year: packed array (1..4) of chari
mpct, fpct: array (1..io) of integer
end:
dentifiers = file of datd:
var
temp: data;
Info: ldentifiers;
fllename: packed array (1..10) of char:
procedure skipliness
var
: integer:
beain
for 1 := 1 to 10 00
writeln
end {skiplines} ;
procedure access-f1le-data;
var
1: integer;
begin
writeln'tenter file nove: ,)
readin(filename);
reset(info, filename);
temp,country : = infon.country;
get(info): (
get(info):
cemp.year := infon.yea
for i iz 1 to 10 do
beqin
temp.mpct(1):= inf0*.mpet[1];
qet(info)
end;
for 1:= 1 to 10 तo
begin
temp.fpet(1) := info^.fpct(1);
get(inf0)
end;
close(1nf0)
end {access-file_data},
procedure putlish_data:
const

```

```

    1: integer:
    beain
skiplines:
writeln(senarator):
writeln:
*ritein(', the followina inforination is⿱⿰㇒一㐄口
writeln(. contained in ', filename, ':');
writeln;
writeln(", ', temr.country):
writein;
"rateln(", , temn.vear):
writein;
write(` rale rercent:'):
for 1:= t to 10 do
write(temp.tect(1): 3):
writel?
writer(' female cercent:');
for 1:= 1 to 10 do
for 1\&:= 1 to 16 do : 3);
writeini

```
    1: integer
    selector: char:
    beain
        writelns
        writeln:
    writeln(. you ray rake any number of changes \({ }^{\circ}\) )
    writein(" fy selectin a
    writelnc
    writeln( \({ }^{\text {w }}\) use the followina set of selectors:')
    writelnf use the tollowino se
writeln!
writeln
    writelri
    writeln( area name......

    writeln( female percent ... fof. \(^{\text {. }}\) )
    writein:
    writeln(e enter the selector, then <ct>, and \({ }^{\circ}\) )
    writeln( you will be prompted to enter the"):
    writeln(" nex data."):
    writeln:
    writeln(" when you have completed the chanaes,")
    writein( enter an \({ }^{\circ} e^{\circ}\) to end the session.");
    writeln:
    write( enter a selector: e)
    readin(selector):
    repeat
            case selector of
                - A', 'a':
                    beain
                    writeln(e enter new area name. \(\left.{ }^{\circ}\right)\);
writelr(e "ise is colurns:
                readinctere.colntry)

                begin
                    writec enter the new year: ,,
                    readinctemp,year
                -mend: me:
                    réain
                    writelnce enter all sixteen male rercent values ...)
                    writeln( \({ }^{\text {nignest age arnuns to lovest: }) \text {; }}\)
                    for \(1:=1\) to 16 do
                    read(temp.mpet(1)):
                    readin

                    beain
                    writeln(e enter all sixteen terale feycent valles ...)
                    writeln( \(n\) nanest acle grours to lowest:'):
                    for \(1:=1\) to 16 do
                    read(temr.fpet(1)):
                    readin
                pnd:
            end (case) :
        witein( make anotner selectinn: \(\quad\) )
        readin(selector)
    intil (selector \(=\) 'E.') or (selectar \(=\) 'e');
    rewrite(info, filename):
    infon.country: \(=\) terre.country
    put(info):
    nfo.year : \(=\) temn.year:
    for \(1:=1\) to 16 do
        begin
            infon. \(\pi\) pet \([1]:=\) tern.rpet[11;
            put(1nfo)
        -ent:
    for \(1:=1\) th 10 do
        or infin
            infon.fpet(1): : terp.fnct(i):
            put(info)
        end:
    lose(info):
    writeln(" new data stored in ', filename)
    witeln:
    oublisn_data
end (make-file_chances)
procedure modify-file_cralce:
    var
    choice: char,
    beain
        writeln;
        riteinc do you want to modify tre data? ....)

        readinchoice):
        if (choice \(=\) 'y') or (cnoice \(=\) 've) then
            make_flle_changes
        mak
            writean( \({ }^{\circ}\) no changes to tne file."
    end (modity-file_choice) ;
begin
            access-flle_data
            publish_datal
            rodify-tile_cnolce
    ent.
3) drawpyramid - The final program accesses the information stored in the external file and produces a
pseudo-graphic on a line printer. The program can produce a single plot, as shown in the BERLIN example, or a double plot of either one region in two time periods or two different regions. The student selects single or double plot and enters the file names. The program takes over from there and produces the output. A double plot of SWEDEN and MEXICO illustrates the age structures of a country with a low rate of growth and one which is high.
```

run getpyr

```
enter file name: sweden.pyr
the following information is contained in sweden.pyr:

\section*{sweden}

1970
male percent: \(1 \begin{array}{llllllllllllllll}2 & 2 & 2 & 3 & 3 & 4 & 4 & 4 & 4 & 3 & 3 & 3 & 4 & 4 & 4\end{array}\)
o you want to modify the data?
if so, enter a 'y'
if not, enter an ' \(n\) ': \(n\)
no changes to the file.
Ready
prograr drampyramid(input, outeut);
(a program to produce a poculation pyramit qrafnie)
const
blonk \(=\cdot \cdot ;\)
type
ddta \(=\)
country: packet array (1..15) of char:
year: Dackec array \(\{1 . .41\) of crar; male, female: array \((i, ., 1 o)\) of inteqer: end)
filename = packed array (1..10) of char:
var
choles: char:
temp: data;
matrix: arrav (1...42, 1..63) of enar:
pyramid: Identiflers:
fllei. tile2: fllenare;
procedure inftializemrray; \{set all array elements to plank\}
var 1, inteaer;
beain \(f:=1\) to 42 ao
for \(1:=1\) th a 300
martixit, 11 := rlanks
procedure riot_cnoice; \{simile plot or superincosen rlot
reain
writeln:

readin(choice)
writeln
ent \{plotachoice) :
procedure enter_file_nare;
negin
 neain
writein( \(e n t e r\) each file nare on a separate line; \({ }^{\circ}\) ):
writeln( enter each file nare on a separate line; ; ;
writelf(..... use ten colurns for enchi note the marker,
writeinc. if you are at the line printer,");
writeinc
writeinc position the writing nead to the \({ }^{\circ}\) );
writeln(
writelnc. lastitine oin the paper before \()\) :
writeln(. enterina <cr> after the second file name.");

readin(filel);
reaoln(filez)
```

    end
    else (sinale flot)
    reain
    writelr(' enter the flle name usina ten columns;')
    writeln(. note tre tarker.......):
    writeln(. it you are at the line printer,")!
    ritelr(. cosition the *ritina read to the"):
    witelr(' last line of the naper before"),
    writein(' enterino <cr>.'):
    writeln;
    writelr(c..: l'):
    reacln(fllei)
    ent
    end {enter-f11f-nare};

```
proceure larels:
    var \(\quad\),,\(k\) integer:
        shorttitle: packed array \(\left(1 . \mathbf{i}^{18}\right)\) of char;
        lonytitle: packed arrav \(\{1 \ldots 3 \mathrm{~b})\) of char:
        agegroups: packed array ( \(1 . . .65\) ) of char:
        menwomen: nacked array \([1, .10]\) of cnar:
    beain
        for \(1:=2\) to 4200
        bealn
        ratrix \([1,1]:=:!;\)
        natrix(1, 63) \(:=\cdot!\cdot\)
        end
        or \(1:=1 \operatorname{tn} 63\) do
            reain

            matr;
        for \(1:=0\) to 3790
        matilx(1, 27): \(=\boldsymbol{l}^{\prime}\);
    for \(1:=7\) to 47 ac .
        matrix(37, 11 : \(={ }^{\circ} \mathbf{- c}^{\prime \prime}\)

            beain
            lonatitle : = porulation Dyranids \(=\) sucerimnosed
            for \(\mathrm{t}:=3\) to 37 do
            matrix end \(\left.^{2}, 1\right):=\) longtitle(1-2)
        end
else
    else
        heain
        shorttitle : \(=\) forulation nyramid',
            for \(1:=3\) to 20 do
            matrix \(\{3,1]:=\) shortititle - 2\(\}\)
        enत:
    enwomen : \(=\) eraleterale'
    for \(y:=13\) to 16 do
        matrix[13, 1\(]:=\) menwomen(4-12):
        for \(y:=38\) to 43 तo
        matrix[13, 1\(]:=\) renwarency - 33):

    Tatrix \((4,55):=e^{\circ}\),
    \(\operatorname{matr} 1 \times(6,53):=\cdot 7 \cdot\) :

    gegroups :=
            - \(70-7405-6960-6455-5950-5445-4940-4435-3930-3425\)
                -2920-2415-1910-14
    \(1:=1 ;\)
\(k:=8 ;\)
    wile \(1<=65 \mathrm{do}\)
        neain
            for \(y:=52\) to bh do
            beain
            inatrix[k, y) \(:=\) ageoroups(1)
            \(1:=1+1\)
            end:
            \(k:=k+2\)
        end:
    matrix(34, b3) \(:=5^{\circ}\) :
    Tatrix(34, 54) \(:=0 \cdot 0 ;\)
    ratrix(34, 55) \(:=90 ;\)
    matrix \([36,53]:=00:\)

    \(\pi\) atrix
1
\(t=7\)
    while \(+<=47\) do
        beain
        motrix[38, 1 ) \(:=\cdots\);

        atrix[38, 27]: \(=0\)
    matrix \(\left(30^{\circ}, 7\right):=010^{\circ}:\)
    matrix \(\left(30^{\circ}, 8\right):=\bullet \circ \cdot\)
    ratrix 39,11\():=080 ;\)
    matrix(39, 15\():=06{ }^{\circ} ;\)
matrix 39,19\():=-40^{\circ} ;\)


    matrix \((39,35):=" 4{ }^{\prime} ;\)
\(\operatorname{matr} 1 \times(39,39):=06 \%\)

    matrix \(\left.139^{\prime}, 4 \mathrm{f}\right):=01^{\circ}:\)
    Totrix(39, 47) \(:=00^{\circ}\)
    ratrix(41, 24): \(:=\) 。
    Tatrix(41, 25)

    Tatrixp(41, 2R) \(={ }^{\prime} e^{\prime} ;\)
Tatrix(41, 29\():=n^{\prime}:\)
    end (larels)?
procedure retrieve.ana_assion_ तata:
```

    var, j, k: integer:
    l, j, k: integer:
    veain
    If (cholce = 'D') or (cholce = "ה") then (double plot)
            beain
            fllearray(1):= f11e1:
            filearray[2] := f1le?;
            for 1:= 1 to 2 do 
                beain
                if 1}=1=1\mathrm{ then
                else
                else
                reset(Dyramit, fllearrayrif):
                remp.country:z pyramido.country,
                get(pyram1d):
            if i =2 then
            for 1 := 1 to 1b do 
                matrix[4, 1 + 17) := temf.country(1)
            else
                for f := 1 to 15 do
                    matrix[4, , + + do := temo.country{j],
        cemp,year := pyramid`,vear:
        get(DYramid);
            for j := 1 to 4 तo
                matrix(5, + + + 17) := temn.vear(])
            else
                forty:=1 to 4 तo 
            for t := 1 to 10 in
                beain
                terp.maje[t] := pyramid*.male(f):
    ```


```

                else
                else
            ford:
        for y:=
                beqin}\mathrm{ terp,ferale(f) := nyramidatamale(i):
                terp.ferale(j)
    ```

```

                matrix{{2* 1 + 4, 27 + 2 * temF.female{{]}<> Dlank then
                matrlx{2 * ) + 4, 27 + 2 * terc.female\f]) := '='
    ```

```

                matrix(2 *
        close(pyramid)
        end
    else (single clot)
    hegin
        reset(ryramit, filel);
        temp.country := pyramid*.country:
        get(pyramid):
        for y := 1 to 15 to
        matrix(4,, ) + 2):= temo.country(j);
        temp.year:= curamtdoyear;
        get(pyramid):
        for t:=1 to 4 तo
        matrix(b, ) + 2):= temp,year(t);
        for }1:=1\mathrm{ to in go'
            beain
            temp,male(1) := nyram1d*.nale(1);
            get(pyramid):
            matrix[2* j + 4, 27 - 2 * tern.maleljl) := **'
        end;
        for fi= 1 to 1t ac
        beain
            terp.temale[t]:= nyra*i*^.ferale[t]:
            qet(Dyramid):
            matrix[2 * ] + 4, >l + 2 * terr.ferale{jl] := '**
            end:
        closec
    ent
    end {retrieve_and_assian_data} ;

```
procedure symbolgexplanation:
    var
        eneck1, eneck 2: datas:
        beain
            wititelf;
            reset (pyramit, fllel);
            creckd.country
            crecki.country
aet(nyramid)
cnecki.year :
            cnecklyear : \(=\)
close(pyratid):
            close(Dyramid): furatita, vear,
            reset(Fyramic):

            aet(Dyramid):
            cneck2.year := fyramia*year;

            if checki.country \(=\) eneck 2 .country then (same area for both plots)
            neain
                    "riteln(: \(\quad\) checkivear: \(\left.4,:-+^{\circ}\right)\) :
                    writeln(. \(\because\) checkivear: \(\left.{ }^{4}, \quad \cdots+{ }^{\circ}\right)!\)

            enc
            else (different areas)
            heain
writel
            writeln( (, enecki.country, ....*);

            writein( \(\quad \therefore=\cdots\) same value for both areas \({ }^{\circ}\) )
            end
        end \{symbol-explanation\};
```

    beain
        for 1:= 1 to 12 Jo
        writelo
    end (skjelines):
    procedure rronllce_ryramiA;
var, y: integer:
seain
tor 1 := 1 to 42 तo
for f:=1 to h3 dn
neair
write(matrix(1, 11);
1t j=03tren
writeln
end;
1f (croice = 'r') or (choice = 'de) then
synrod_explanation:
skifilires
enत {pronuce_pyramid)
beasr
Inftlalize_array:
rlot_crolce;
enter-file-nare;
larels:
retrleve-ard_assion_Aata;
frocucembyrania
end.

```


\section*{Conclusion}

The inclusion of exercises such as this one in social science classes has proven to be valuable in a number of ways. It allows students with little or no programming background to get over their tentativeness about approaching a computer. In addition, I believe that such exposure to computers, however limited, contributes to the overall computer literacy of students. Finally, the experience may spur a student to want to take a course in computer programming or to learn other uses of the computer.

There is absolutely no reason why students in all divisions of the liberal arts setting should not benefit by

```

sweden
mexico -- +
$1='$ : same value for both areas

```
the opportunities available in the field of computer science.

\section*{Notes}
1. The Geography Program Exchange assists universities and other non-profit organizations with the interchange of computer software which relates to problems of a geographic nature. The address is:

\section*{Geography Program Exchange \\ Department of Geography \\ Michigan State University \\ East Lansing, Michigan 45824}
2. One of only a few books written on the general topic of computer applications in geography is Paul M. Mather, Computers in Geography: A Practical Approach (Oxford: Basic Blackwell, 1976); it contains four chapters, one of which is an introduction to the FORTRAN language.
3. Mark S. Mormonier, Computer-Assisted Cartography: Principles and Prospects (Englewood Cliffs, New Jersey: Prentice-Hall, Inc., 1982), p. 22.
4. John L. Kirkley, Editor, ''Our Industry Could Lead a Liberal Arts Renaissance,'" Datamation, March, 1983, p. 29.
5. Brian J. L. Berry, et al, The Geography of Economic Systems (Englewood Cliffs, New Jersey: Pren-tice-Hall, Inc., 1976), p. 36.
6. The programs in this paper were prepared using Oregon Software Pascal, Version 2.0, and run on a DEC PDP 11/70.

PUG

\title{
Path Pascal A Language for Concurrent Algorithms
}

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}

\section*{1. Introduction}

This paper is intended to provide an overview of the Path Pascal programming language. Rather than introduce the language by studying its definition, the approach taken here is to explore a moderately complex example. While any detailed understanding of Path Pascal must be based upon the formal definition of the language, this paper will present the most important concepts embodied in the language.

After a brief history of the development of Path Pascal, the problem to be solved by the example program will be presented. Using this example, three major concepts of Path Pascal will be explored. With these major concepts, the operation of the program can be understood. Finally, a summary of the present status and anticipated future of Path Pascal are discussed.

\section*{2. Background}

Path Pascal was originally developed at the University of Illinois in 1978 by Dr. Roy Campbell. Details of the original Path Pascal compiler project are available from the University of Illinois as a series of Research Reports. In addition, the definition of the language has appeared in SIGPLAN Notices [1].

The University of Illinois implementation of Path Pascal was for the LSI-11/23 processor. Path Pascal has now been implemented on a variety of machines including the M68000 (by NASA-LaRC), AMAC-80 (by Martin-Marietta) and VAX-11/780 (by NASA-Goddard). This report is based upon experience gained with Path Pascal as part of NASA Contract NAS1-16985 during 1982 using the M68000-based Path Pascal compiler developed by Dr. Ed Foudriat at NASA's Langley Research Center.

Path Pascal is based upon 'Path Expressions'' first described by Campbell and Habermann in 1974 [2]. The key concept is that coordination among a collection of concurrent processes should be expressed in a language designed especially for that purpose. "Path Expressions" conveniently and succinctly specify the central concepts of "mutual exclusion" (protecting "critical sections'" of code) and of "synchronization" (waiting for information to be computed by other processes). In Path Pascal, the primary unit for mutual exclusion and synchronization is the subroutine, allowing the use symbolic names in the Path Expressions.

In Path Pascal, 'counting semaphores'" are used to implement both mutual exclusion and synchronization. By specifying the Path Expression prior to the subroutines that it controls, the compiler can generate appropriate initialization, P-operation and V-operation at the beginning and ending of each subroutine.

\subsection*{2.1 The Island}

Before discussing the example Path Pascal program, it may be useful to understand that this program is a simple event-driven simulation. The simulation involves an island and its inhabitants.

The island of this program is a very special island. It consists of a \(25 \times 17\) grid, each element of which can either be empty (displayed as a blank), contain a wolf (designated by a ' \(W\) '), or contain a rabbit (designated by a ' \(R\) '). Initially, there are 17 wolves (all in column 10 ) and 17 rabbits (all in column 16).

Each wolf begins with a user-specified "energy". This energy is used on an "annual' ' basis to remain alive, looking for rabbits to eat or other wolves with which to mate. Each 'year'' the wolf looks around his position on the grid, determining if there are any rabbits or wolves in his neighborhood. If there are any rabbits, the wolf's energy is increased by eating them. If, on the other hand, there are too many wolves in the neighborhood, the wolf loses excess energy due to overcrowding. Only if there are a reasonable number of neighbor wolves and this wolf is "fertile" does the wolf attempt to produce an offspring. If the wolf's energy is reduced to zero, it dies.

Each rabbit also begins with a user-specified "energy''. This energy is affected in a manner similar to a wolf, except that a rabbit is considered to have been 'eaten'" if there are any wolves in its immediate neighborhood and that rabbits gain energy by overcrowding rather than losing it.

Finally, the user may wish to "repopulate" the island, assigning new energy and lifetime specifications. This is done by pressing any key on the keyboard.

\section*{3. Path Pascal Constructs}

The example program, ISLAND, is primarily written in standard Pascal. It makes use of only three new constructs - OBJECTs, PROCESSes and Wait-for-Son processing.

\subsection*{3.1. OBJECTs}

Of the several extensions to standard Pascal, the OBJECT construct is the most important to understanding Path Pascal. An OBJECT is a Path Pascal TYPE with several properties similar to a RECORD. As with a RECORD, each variable of an OBJECT TYPE allocates stack space and NEW of a pointer to an OBJECT TYPE allocates heap space. The space required for a OBJECT TYPE is that for the semaphores implied by the OBJECT's Path Expression and for any variables explicitly declared within the OBJECT. Unlike RECORDs, OBJECTs contain subroutines (PROCE-

DURESs, FUNCTIONs and PROCESSes). Those subroutines that are included in the Path Expression are termed "ENTRY"' routines and may be accessed from outside of the OBJECT using the RECORD-like notation "object.entry(parameters)" (e.g., line 436 or line 444).

Lines 16-52, 71-95, 96-183 and 186-200 indicate four common kinds of OBJECTs. These examples will be discussed in the following sections.

\subsection*{3.1.1. OBJECT CRTOBJ (lines 16-52)}

This OBJECT is an example of an "interprocess buffer" OBJECT. Since processes run concurrently (see Section 4), they must be synchronized in order to transfer information. Unlike ADA in which processes must "rendezvous", Path Pascal facilitates the concept of "interprocess buffers" that contain data to be transferred from one process to another. This allows the "sending" process to continue execution after generating the information for the other process.

The Path Expression on lines 19-21 has both mutual exclusion and synchronization expressions. the first two expressions simply state that the operations PUSH and POP are atomic (only one PUSH at a time and only one POP at a time). The last expression specifies the synchronization between PUSH and POP. It states that a call to POP may not proceed until a call to PUSH has completed and, furthermore, at most CRTSZ calls to PUSH can be honored before at least one call to POP occurs.

This last implication of the Path Expression, that at most CRTSZ calls to PUSH can proceed without at least one call to POP, is the key to understanding the data structures (lines 24-27) and code (lines 29-51) of CRTOBJ. Since at most CRTSZ calls to PUSH can occur without a call to POP, all that is required is space for CRTSZ "messages". Since the "interprocess message" in this case is just a character, BUF is simply an ARRAY of CRTSZ characters. INPTR specifies where PUSH is to put its character, and OUTPTR specifies where POP is to get its character. This code works because the Path Expression controls access to the routines PUSH and POP, and, therefore, controls access to the BUF ARRAY.

\subsection*{3.1.2. OBJECT CREATOR (lines 71-95)}

This OBJECT is also an interprocess buffer; however, the buffer has only a single entry (the VARs XX, YY and EE). This form of the interprocess buffer is very similar to the ADA "rendezvous".

\subsection*{3.1.3. OBJECT SCREEN (lines 96-183)}

This OBJECT is used to control access to the INFO ARRAY, the inmemory representation of the island. The Path Expression on line 99 simply states that one and only one of the allowed operations may be progress at any given instant.

The routines in this OBJECT include SETUP (for reinitialization), KILL (for termination), LOOK (for examining the "neighborhood" of a wolf or rabbit), ASSIGN (for direct control of INFO), CHANGE (a test-and-set operation) and DONE (a set-and-test opera-
tion). The routine WRITES (update terminal screen) is available only within this OBJECT.

\subsection*{3.1.4. OBJECT SHUTUP (lines 186-200)}

This OBJECT is used to synchronize the termination of the simulation. Since Path Pascal does not allow "preemptive termination" of process (see Section 4), care must be taken when writing processes that must eventually terminate.

This OBJECT acts essentially as a binary semaphore. The call to SHUTUP. WAIT on line 376 will cause the calling process (SHUTDOWN) to suspend operation until the call to SHUTUP.SIGNAL is made on line 364.

\subsection*{3.2. PROCESSes}

The second major addition of Path Pascal to standard Pascal is the PROCESS. Conceptually, a PROCESS is a PROCEDURE that, after it is called, executes in parallel with its caller. In Path Pascal, all processes that are not waiting due to a Path Expression, a DOIO (see below) or a DELAY are competing for the hardware processor(s). Also note that each call to a PROCESS creates a new process, as in lines 311-317.

PROCESSes in Path Pascal can either be normal or INTERRUPT PROCESSes. An INTERRUPT PROCESS has two special attributes not associated with normal PROCESSes. The PRIORITY and VECTOR information are used to control the interrupt hardware such that the DOIO statement (lines 213-237) acts as a "wait-for-interrupt". In addition, as shown on lines 211, 216, 234, 239, it is sometimes necessary to enter "supervisor state" in order to access device controllers.

\subsection*{3.2.1. INTERRUPT PROCESS DLVJIN (lines 201-221)}

This process is an "infinite loop", waiting for an interrupt from the terminal input hardware. When such an input occurs, the character is forwarded to the appropriate interprocess buffer.

\subsection*{3.2.2. INTERRUPT PROCESS DLVJOUT (lines 224-241)}

This process is also an "infinite loop". It waits for a character to be placed into the appropriate interprocess buffer, transfer the character to the terminal, and waits for the completion interrupt.

\subsection*{3.2.3. PROCESS WOLF (lines 242-270)}

This process corresponds to a wolf in the simulation; it is called when a wolf is to be created. The process is a loop corresponding to the lifetime of the wolf. In the loop, the ENERGY of the wolf (a local variable) is constantly updated until it is reduced to zero and the wolf dies.

\subsection*{3.2.4. PROCESS RABBIT (lines 273-302)}

This PROCESS is similar to PROCESS WOLF, except that it corresponds to a rabbit. As with the wolf, the ENERGY of the rabbit is constantly updated until it is reduced to zero and the rabbit dies.

\subsection*{3.2.5. PROCESS SHUTDOWN (lines 366-380)}

This PROCESS is used to wait for input from the user (any input will do). When this input occurs, it is necessary to notify the screen monitor (line 374) and the main program (line 375). However, at thia point it is necessary to wait for the main program (actually, PROCEDURE PROCREATE) to complete its processing (line 376). Finally, SHUTDOWN 'absorbs'" any extra attempts to create rabbits or wolves. This is completed when the special message having an ENERGY of zero is encountered, and SHUTDOWN is terminated.

\subsection*{3.3. Wait-for-Sons Processing}

When a PROGRAM, PROCEDURE, FUNCTION or PROCESS calls a PROCESS, it is necessary that this "son'" process terminate before the "father" can terminate. This is logically necessary due to the scope rules of Path Pascal. Furthermore, this 'wait-for-sons" processing is a useful tool for coordinating the termination of a system.

Except for ' wait-for-sons'" processing, there is no reason that the code in PROCEDURE PROCREATE could not be part of the main program. Note, however, that all of the WOLF and RABBIT processes are initiated by PROCREATE. Hence, PROCREATE cannot continue until all of these processes have terminated. This fact is critical to the coordination between PROCREATE and SHUTDOWN when the simulation is being terminated.

\section*{4. PROGRAM ISLAND}

Having looked at the special features of Path Pascal that are used by this program, it is now possible to step through a typical execution of the program.

The main program begins (lines \(430-433\) ) by allocating heap-space for CRTIBUF and CRTOBUF, by initiating the input/output processes DLJVIN and DLVJOUT, and associating DLJVIN with CRTIBUF and DLJVOUT with CRTOBUF.

The driving loop of the program (lines 435-445) clears the terminal's screen (using DEC-VT52 protocol), prompts the user for parameters (PARAMS), reinitializes the simulation (SCREEN.SETUP), initiates a process to look for terminal input (SHUTDOWN), and calls PROCEDURE PROCREATE.

PROCEDURE PROCREATE (lines 303-365) begins by initiating 17 wolves and 17 rabbits. It then enters a loop waiting for requests for creation. When such a request occurs, it is first tested to see if it was generated by SHUTDOWN, indicating that termination should begin. If this is not a SHUTDOWN request, it is a request for the creation of a wolf (ENERGY \(>0\) ) or a rabbit (ENERGY<0). Each direction (UP, DOWN, LEFT and RIGHT) is tested to see if it is available. If all directions are occupied, creation is not possible. If a free position on the island is found, the SCREEN.CHANGE call updates the simulation and the rabbit or wolf is created (lines 358-361). When the special SHUTDOWN request is encountered, PROCREATE signals SHUTDOWN that it has completed processing, and waits for all of the wolf and rabbit processes to terminate.

Once all of the wolf and rabbit processes terminate, PROCREATE returns to the main program (line
444). The main program now signals the SHUTDOWN process that no more requests for creation will be generated, and the master control loop iterates.

\section*{5. Summary and Conclusions}

While a single example cannot cover all of the constructs and uses of those constructs, the ISLAND program is representative of the important capabilities that Path Pascal has that are not found in standard Pascal. These capabilities include multiple processes (PROCESS), interprocess coordination (OBJECT), and process termination coordination (Wait-for-Son).

Having programmed in Path Pascal for several months, it is clear that these new capabilities are useful. Many of the PROCESSes and OBJECTs that have been written have been found to be highly reuseable since they "encapsulate" and entire concept or function within the program. However, it is equally clear that these new capabilities do not "solve" the concurrent programming problem. Developing the Path Expressions is a tedious, error-prone undertaking. Nonetheless, once a Path Expression is finally 'correct', it is usually clear to anyone reading the code exactly what will occur when the program is executed.

One of the goals of the current research with Path Pascal is to identify various "prototype"' Path Expressions. The 'interprocess buffer"' is a good example. If a few such prototypes can be found to be sufficient for most situations, a "macro OBJECT' facility might be added to Path Pascal to make these prototypes readily available to the average programmer.


```

            RECIN
            REAT(IR.(REATE( X,Y, ENER(;Y)
                IF ER. R(`R()WI) THEN (* RABRITS I,OVE A (`ROWD *)
                F.NF:RGY : F:NE:RI;Y+2•RANNIMAI;
    ```
\(\qquad\)
```

            ENF:RI;Y :-
        F ENERGY & RMAX THEN
            ENERGY : RMAX:
        Y :- ENERGY RMAX
        IF DL,Y&- ENERGY THEN
        DIY :- -ENERGY:
        DELAY(DLY);
    (NTIL. SCREEN.DONE (X,Y,ENERGY); (* DECIDF. WHETHER STILL "ALIVE" *)
    END: (* PROCESS RABBIT *)
PROC
X \
DIRS
GIN
FOR Y :- 1 TO YMAX 1 DO (* INITIALIZE SCREEN *
REGIN
SCREEN.ASSIGN(XWOLF,Y,WINIT); (* A COIUMN OF WOL,VES ;,
WOLF( XWOLF,Y,WINIT);
SCREEN.ASSIGN(XRABBIT,Y,RINIT); (, A COLIIMN OF RABRITS.
RARBIT(XRABRIT,Y,RINIT);
ENI;
STR :-- UP: (- CREATION DIRECTION *)
RFPFAT \&* CREATE OFFSPRING AS REQUIRED *)
CRFATOR.STARTUP(X,Y,ENERGY); (* OBTAIN OPERATION REQUEST *)
IF ENERGY >0 THEN (, RFQUEST FOR CREATION),
BEGIN
DIRS : 4: (* TRY ALI FOUR DIRFCTIONS *)
RFPEAT DIR-RIGHT THEN (. (ONSIDER NEXT DIRE(TION .
DIR-RIGHT THEN
EI.:FE DIR : SOC(CDIR)
(MASF: DIR OF
IF SCRFEN.(HANGE(X,Y 1, ENERGY) THEN
MEGIN
DIR:; : O: (: SET FLAG \& )
END:
if S(CREEN.(HANGE(X,Y+1) ENERGY) THEN
mplilN
WIRS: 0): (: SET FlAG \&
END:
LEFF : IF X.I THEN
IF' SCREEN.(HANGE:(X ].Y, ENER(;Y) THEN
BFC;IN
HIR:; : 0: ( OET FLAG; •
\
K:IHT : IF X+2.XMAX THEN
IF SiCRFEN (HANCFF: X + I Y Y ENE:R(iY) THEN
RE,IIN

```

```

                X
                    EN|!:
                WIRC: H:R:B 1:
                INTII WIR:; O:
                    IF ENFRGY.)(. IF 0. (ANNO)T [M)(REATION .)
                    F ENER(;Y - O THEN , ' 'REATE NEW RARBIT ',
                        RARBIT(X,Y,FNEK.;Y,
                    FH:MF
                ENI):
    ```

```

    SHIMTUP.:II;NAL:: ( SHITIOWN HANDIF:S EXTRANF:OUS (RFATION RFQUF:ST:;
    ENL: (. FROCEDOIRE PRIM-REATE: )

```
```

PROCESS SHUTDOWN; (* HANDLE TERMINATTION *)
CH : CHAR; (* INPUT CHARACTER *)
CH : CHAR; (* INPUT CHARACTER *)
Y ENFRGY : INTEGER;
BEGIN
CRTIBUF^.POP(CH); (* WAIT FOR TERMINAL INPUT *)
SCRFEN,KILLL;
(* NOTIFY SCREEN MONITOR *)
REATOR.CREATE(0,0,0); (. TERMINATE PROCEATION*)
HIJTIP. WAIT
REPEAT (* HANILLE ANY EXTRANEOUS ATTEMPTS AT CREATION *)
CRFAT(IR STARTUP(X,Y. ENERGY):
INTIT. FNERGY-O;
FNN): (= PROCFSS SHITTD)WN *)
PROCEINOFF. PARAMS: (. PROMPT USER FOR INPUTT *)
TYPE
STRIN(;G, PA(KEI) ARRAY ! 1..t1 OF (HAR
pRIMFIDILR: FARAMIN(X,Y:INTE:GER; M:GG:STRIN(GG; VAR VAI,:INTEGER);
var
I : INTF:GFR
CH: 'HAR
RFC:IN

```

```

        RTMRUF`.FISH(CHR(FSC))
        *TORIIF', PULOH('Y'):
        (RTCORIFF- FUSH(CHR(Y+ 11));
    ```

```

        (RTORUF'.PUSH(MSGGI1);
        VAI : ""; (. RFAI) ANE DECODPF VALUF. ,)
        REFFAT IRTIRIJF".POP(CH):
            *K'IRIFP`.POP(CH): 
                MBFOIN
                    VAI, VAI.-1O+(ORI)
                FND
            HN
                IF (H.,CHR(CR) THEN (. UNRF(OMNIZFI) *)
                (PTORUF*. PU:AH(CHR(BEI. i)
        |N:II, VA,.G+FH) UR :(H (HR(CR)):
    FNN!: i · FR:, FlumF: :ARAmin *)
    MFOIN
FARAMIN/ XMAX+\&.1.'WINIT:', WINIT); (' INITIAI, ENF.RGY FOR WOLVES ·)
ARAMIN(XMAX+G.?, RINIT:`, RINIT'i:,' INITIAI, FNFRGY FOK RABHITS, ')

```

```

    WANNIAL: W:N:'MIV WIRS: : WUTF VAIHF:' ARE POSITIVE ·)
    WFFRTILF: |WENNIAL::
    wi/k\WI: : \WiN:T
    MAX : -WINIT
    KINIT: KINIT: (' RABRIT VAMNF:S ARE. NEGATIVFE .)
    FANNIAI: : RINIT IIV RYRS:;
    HFF:RTIIF: - 2-RANNIAI::
    (c(R(WI): I.RINIT
    max : zPGINIT:
    ENI, ; procemitrf: farams .,
KHIIN - IFH:MAM :IIAN
NFW(` KTIRUF': (, 'REATE: INP(OT BIIFYE:K ))
MFW(iRTIM(IF):

```

```

    DIV.JOIFT('RTORIIF): (. :TARTUF' OUTPUTT PROCESSS ,
    RFPEAT
        `RTOBUF`. PH:AH(CHR(FEG`)):
        'RTOBUF' PISH('H')
        CRTOBUF'.PUSH(CHR(ESC)): (. VT52 CLEAR SCREEN OPERATION *)
        "RT(IRUF'. PuNG('J');
        PARAMS:
        SCREEN .SETUP;
        SHUPTDWN:
        PROCREATE
        M,
        CRFATOR (* RETURNS WHEN ALI, TASKS COMPLETED *)
        CRFATOR.CREATE(0.0.0); (* TERMINATE SHUTDOWN*)
    ```
    UNTII. FALSE:
ND. (* PROGRAM ISLAND *)

\title{
An Introduction to Modula-2 for Pascal Programmers
}

\author{
By Lee Jacobson and Bebo White \\ Jacobson, White, \& Associates \\ San Francisco, CA
}

\section*{THE BACKGROUND AND HISTORY OF MODULA-2}

Modula-2 (like Pascal) was developed at the ETHZurich under the direction of Niklaus Wirth (Institut fur Informatik). Its development grew largely from a practical need for a general purpose, efficiently implementable systems programming language. The first production use of Modula-2 occurred in 1981. Dr. Wirth's book, 'Programming in Modula-2' was published by Springer-Verlag in 1982.

It is virtually impossible to examine Modula-2 without recognizing its roots in Pascal. In its original design, Pascal was intended to be a language suitable for teaching programming as a systematic discipline based on certain fundamental concepts clearly and naturally reflected within it. These concepts were largely centered around stepwise refinement of problem solutions and structured programming.

Inasmuch as Pascal is basically an academic language, its widespread use for a variety of applications has clearly exceeded its design intention. Hence, many extensions to the original Pascal definition have been designed. Likewise, it has attracted as many critics as it has disciples.

Modula-2 has assumed all of the positive features of Pascal, and has attempted to address its commonly recognized shortcomings. The result is a structured, modular, portable, readable, efficient, machine independent, flexible language.

This paper will address the primary differences between Modula-2 and Pascal with particular emphasis on some of those features which the authors consider quite significant. Programming examples will be given in both Modula-2 and Pascal.

\section*{MODULA-2'S DIFFERENCES FROM PASCAL}

\section*{The Role of Modules in Modula-2}

Modules are the most important feature distinguishing Modula-2 from Pascal. Relying heavily upon the concepts of scope and block, modules address the problem, usually found in large programs, of separating visibility from existence. In block-structured languages, the range in which an object (e.g. a variable or procedure) is known is called the object's scope, and therefore, defines its visibility. However, an object's visibility also binds its existence, in that objects are created when the block in which they reside is entered and destroyed when the block is exited. It should be possible to declare variables that maintain their values, but are visible only in a few parts of a program. Concurrently, there is also a need for closer control of visibility. A procedure should not be able to access every
object declared outside of it when it only needs to access a few of them.

Syntactically, modules closely resemble procedures, but they have different rules about visibility and the existence of their locally declared objects. Consider the following declarations:
```

PROCEDURE Outside;
VAR x,y,z: INTEGER;
MODULE Mod;
IMPORT K;
EXPORT a,P1;
VAR a,b,c: INTEGER;
PROCEDURE P1;
BEGIN
a := a + 1;
x := a;
END P1;
END Mod;
END Outside;
PROCEDURE Outside; VAR $x, y, z:$ INTEGER;
(* no module here *) a,b,c:INTEGER;
PROCEDURE P1;
BEGIN
$a:=a+1 ;$
$x:=a ;$
END; (*P1*)
END Outside;
END; (* Outside *)

```

The only syntactic difference between the module Mod and a normal Pascal procedure declaration are the reserved word beginning the declaration (MODULE instead of PROCEDURE) and the presence of IMPORT and EXPORT declarations following the module heading.

The semantic differences are more interesting. The objects declared within \(\operatorname{Mod}(a, b, c)\) exist at the same time as the variables \(\mathrm{x}, \mathrm{y}\), and z , and remain so as long as Outside is active. The objects named in Mod's IMPORT list are the only externally declared objects visible within Mod (x but not y or z). The objects declared in Mod's EXPORT list are the only locally declared objects visible outside Mod. Thus, a and P1 are accessible from Outside, but b and c remain hidden inside Mod.

Specifically, a module can be thought of as a syntactically opaque wall protecting its enclosed objects, be they variables or procedures. The export list names identifiers defined inside the module that are also to be visible outside. The import list names the identifier defined outside the module that is visible inside. Generally, the rules for modules are:
1. Locally declared objects exist as long as the enclosing procedure remains activated;
2. Locally declared objects are visible inside the module and if they appear in the module's export list, they are also visible outside;
3. Objects declared outside of the module are visible inside only if they appear in the module's import list;

The following example demonstrates the essence of modularity:
```

MODULE MainProgram;
PROGRAM MainProgram;
VAR SEed : INTEGER;

```
```

MODULE RandomNumbers

```
MODULE RandomNumbers
    IMPORT TimeOfDay;
    IMPORT TimeOfDay;
    EXPORT Random;
    EXPORT Random;
    CONST Modulus = 2345;
    CONST Modulus = 2345;
        Increment = 7227
        Increment = 7227
    VAR Seed : INTEGER;
    VAR Seed : INTEGER;
        FUNCTION Random : INTEGER;
        FUNCTION Random : INTEGER;
        CONST Modulus = 2345;
        CONST Modulus = 2345;
    PROCEDURE Random() : INTEGER
    PROCEDURE Random() : INTEGER
    BEGIY
    BEGIY
        Seed := (Seed+Increment)
        Seed := (Seed+Increment)
                MOD Modulus:
                MOD Modulus:
        RETURN Seed
        RETURN Seed
    END Random;
    END Random;
    BEGIN (* RandomNumber *)
    BEGIN (* RandomNumber *)
    Seed := TimeOfDay;
    Seed := TimeOfDay;
END RandomNumber;
END RandomNumber;
BEGIN (* MainProgram *)
BEGIN (* MainProgram *)
    WriteInt(Random(), 7);
    WriteInt(Random(), 7);
    N (* MainProgram *
    N (* MainProgram *
    Seed := TimeOfDay;
    Seed := TimeOfDay;
    Writeln(Random, 7);
    Writeln(Random, 7);
END. (* MainProgram *)
```

END. (* MainProgram *)

```

The random number generator in these examples uses a seed variable to generate the next random number. Thus, the seed must maintain its value across function calls. The program on the right shows the classical Pascal solution. Notice that Seed's declaration is at the top of the program, while its initialization is forced to the bottom. Two obvious disadvantages arise from the scattering of Seed across the face of the program:
1. Its occurences become hard to find, especially in a large program;
2. It becomes accessible to every other procedure in the program even though it is used only by Random;

The example on the left demonstrates the usefulness of the module structure. The only object visible to the ouside world is the procedure Random, while all objects pertaining to the random number generator are contained in one place. Note that the module RandomNumber contains both declarations and a statement part. Module bodies are the (optional) outermost statement parts of module declarations and serve to initialize a module's variables. Although subject to the module's restrictive visibility rules, module bodies conceptually belong to the enclosing procedure rather than the modules themselves. Therefore, module bodies are automatically executed when the enclosing procedure is called.

\section*{Relaxed Declaration Order}

New Pascal users are often frustrated and confused by the enforced declaration and definition block structure required within the program skeleton. Despite the emphasis on modules, blocks still play an important part in Modula-2: implementation modules, program modules, internal modules, and procedures are all declared as blocks. Differences from Pascal include relaxed order of declarations, termination of all blocks by a procedure or module identifier, and the optional nature of block bodies.

Pascal imposes a strict order on the declaration of objects; within any given block, labels must be declared
before constants, constants before types, and so on. Modula-2 eliminates this restriction - declarations can appear in any order. Programs containing a large number of declarations are easier to read and understand when related declarations are grouped together (regardless of their kind).

The following is an example of relaxed declaration order:
```

MODULE Xlator;
CONST MaxsSym = 1024;
TYPE SymBuffer = ARRAY[1..MaxSym] OF CHAR;
VAR SymBuff1, SymBuff2: SymBuffer;
CONST MaxCode = 512;
TYPE CodeBuffer = ARRAYII..MaxCodel OF BYTE;
VAR CodeBuff: CodeBuffer;
END Xlator

```

This example easily demonstrates how various related declarations may be placed together in a Modula2 program, whereas in a Pascal program they may be scattered due to strict block ordering. Relaxed declaration order not only improves readability but enables a logical ordering which may be very important in large programs.

\section*{GOTO-less Programming In Modula-2}

Inasmuch as structured programming is often equated with elimination of the use of unconditional transfers, Pascal was designed to de-emphasize use of the GOTO statement. Still the GOTO statement and the LABEL 'type' were supported to allow programming cases where the Pascal logical structures were insufficient. This meant that a GOTO statement was available for use in a situation which would otherwise have forced restructuring of the program logic.

For example, consider the following two program segments:
```

Remainder := Alpha MOD Beta;
WHILE Remainder <> O DO
WHILE R
EGIN
Alpha:= Beta;
Beta:= Remainder;
Remainder := Alpha MOD Beta
END;
10: Remainder := Alpha MOD Beta;
IF Remainder = O THEN
GOTO :20;
Alpha:= Beta;
Beta := Remainder;
GOTO 10;
20:

```

The example on the left avoids use of a GOTO by duplicating an operation. The example on the right, while using GOTOs is actually more explicit.

Modula-2 does not support Pascal GOTO and LABEL. Instead it provides transfer mechanisms for uses under particular controlled circumstances. One of these mechanisms is the EXIT statement which permits premature exiting of a loop. The following is a program segment in Modula-2 performing the same operation:
```

LOOP
Remainder := Alpha MOD Beta;
IF Remainder := 0 THEN EXIT;
Alpha:= Beta
Beta := Remainder
END;

```

This example also illustrates the Modula-2 LOOP construct which operates as a Loop-Forever structure. When the EXIT statement is executed, program control
will transfer to the statement following the END statement which terminates the range of the LOOP.

Additional examples of unconditional transfers supported by Modula-2 include the RETURN statement which is used to prematurely exit a procedure, and the HALT standard procedure which terminates the current program.

\section*{Dynamic Array Parameters}

Another important distinction between Modula-2 and Pascal involves the capability to declare dynamic array parameters. Modula-2 allows formal parameter types of the form:

\section*{ARRAY OF T}
where T is an arbitrary base type. Note that the array bounds are omitted defining a dynamic array type which is compatible with all (one dimensional) arrays having the same base type T .

The ramifications of this feature are widespread. through it, procedures are able to pass to other procedures (functions, etc.) arrays of unspecified size. (Index checking is accomplished by means of a new standard procedure HIGH).

Perhaps the most important way in which dynamic array perameters may be used is in the area of string processing. This feature lifts the rigid Pascal restriction concerning the value assignment and comparison of string variables. No longer is it necessary that operations may only be performed on strings which have the same length.

\section*{Separate Compilation}

Separate compilation is allowed by the Modula-2 compiler through the use of the compilation unit. Mod-ula-2 programs are constructed from two kinds of compilation units: program modules and library modules. Program modules are single compilation units and their compiled forms constitute executable programs. They are analogous to standard Pascal programs.

Library modules are a different animal and form the basis for the Modula-2 library. They are divided into a definition module and an implementation module. Definition modules contain declarations of the objects which are exported to other compilation units. Implementation modules contain the code implementing the library module. Both always exist as a pair and are related by being declared with the same module identifier.

To understand the rationale behind dividing a library module into separate definition and implementation modules, consider the design and development of a large software system, such as an operating system. The first step in designing such a system is to identify major subsystems and design interfaces through which the subsystems communicate. Once this is done, actual development of the subsystems can proceed, with each programmer responsible for developing one (or more) of the subsystems.

Now consider the project requirements in terms of Modula-2's separate compilation facilities. Subsystems will most likely be composed of one or more compilation units. Defining and maintaining consistent interfaces is of critical importance in ensuring error-free
communication between subsystems. During the design stage, however, the subsystems themselves do not yet exist. They are known only by their interfaces.

The concept of a subsystem interface corresponds to the definition module construct. Thus, interfaces can be defined as a set of definition modules before subsystem development (i.e., design and coding of the implementation modules) begins. These modules are distributed to all members of the programming group, and it is through these modules that inter-subsystem is defined. Interface consistency is automatically enforced by the compiler.

\section*{Modula-2 Libraries}

The library is a collection of separately compiled modules that forms an essential part of most Modula-2 implementations. It typically contains the following kinds of modules:
1. Low-level system modules which provide access to local system resources;
2. Standard utility modules which provide a consistent system environment across all Modula-2 implementations;
3. General-purpose modules which provide useful operations to many programs;
4. Special-purpose modules which form part of a single program;

The library is stored on one or more disk files containing compiled forms of the library module's compilation units. The library is accessed by both the compiler and the program loader - the former reads the compiled definition modules while compiling and the latter loads the compiled implementation modules when executing the program that imports library modules.

A dependency arises between library modules and the modules that import them. Consider the example of a single library module. The compiler must reference the module's symbol file (a compiled definition module) in order to compile the implementation module. Therefore, the definition module must be compiled first. Once an implementation module has been compiled, its object file is tied to the current symbol file, as the object code is based on procedure and data offsets obtained from the symbol file. Similarly, when a program imports a library module, it is assumed that the symbol file offsets are accurate reflections of the corresponding object file.

The Modula-2 language contains no standard procedures for I/O, memory allocation, or process scheduling. Instead, these facilities are provided by standard utility modules stored in the library. Standard utility modules are expected to be available in every Modula2 implementation. Thus, by using only standard modules, Modula- 2 programs become portable across all implementations.

The advantages of expressing commonly-used routines as library modules (rather than part of the language) include a smaller compiler, smaller run-time system, and the ability to define alternative facilities when the standard facilities prove insufficient. Disadvantages include the need to explicitly import and bind
library modules, and occasionally a less flexible syntax imposed by expressing standard routines as library modules (as opposed to their being handled specially by the compiler).

\section*{CONCLUSION}

The examples cited above can only provide a clue as to the power and flexibility of the Modula-2 language. It is the hope of the authors that they can pique significant curiosity and interest into this amazing new programming tool.

\section*{REFERENCES}
1. Niklaus Wirth, Programming in Modula-2, Sprin-ger-Verlag, 1982
2. Niklaus Wirth, MODULA-2, ETH Institut fur Informatik Report No. 36, reprinted by Volition Systems, Del Mar, CA, 1980
3. Rich Gleaves, Modula II User's Manual, Volition Systems, 1982
4. Roger Sumner and Rich Gleaves, Modula-2 - A Solution to Pascal's Problems, Volition Systems, 1982

PUG


Data Structures Using Pascal
by A.M. Tenenbaum and M.J. Augenstein
Prentice-Hall, 1981
This book is intended as a text for a first course in data structures that is also a second course in programming. It presents all of the major data structures including stacks, queues, lists, trees, and graphs and describes recursion, list processing, sorting, and searching. An appendix provides a brief tutorial on Pascal. The emphasis is on practical techniques as opposed to theoretical concepts. All algorithms and examples are presented in Pascal.

This book is excellent both for students and for practicing programmers who want to learn how to apply algorithms and data structures, whether or not they use Pascal. However it would not be appropriate, nor was it intended to be, for those merely wishing to learn Pascal.

The authors employ several pedagogical techniques which others would do well to emulate. First numerous examples and sample programs are presented; the authors do not merely rely on textual explanations. In spite of this, there are very few typographical or al-
gorithmic errors as so often is the case with multiple figures. Second, the same basic figure is repeated revaral times with each version successively updated to show the intermediate results of an algorithm. For example, an array is listed after each pass of a sorting algorithm so the reader can follow how the sort progresses. Third, algorithms are often presented as a combination of Pascal and pseudocode, thus highlighting the key points and not confusing the reader with such irrelevancies as initialization or I/O. Fourth, algorithms are presented several times with each new version a refinement of the previous one.

My only criticisms would be that some of the algorithms could be simplified, frequently by more appropriate tests in "while" statements; and more use should be made of enumeration types - certain aldorithms had a Fortran ring to them. However these are nitpicks: the book is excellent and is highly reconmended to all PUG'ers.

Arthur Salwin 1405 Homeric Ct. McLean, VA 22101


\section*{SOFTWARE BUILDING BLOCKS, INC. ANNOUNCES PASCAL COMPILER FOR THE IBM PERSONAL COMPUTER \({ }^{\circledR}\)}

ITHACA, NY - A new company, Software Building Blocks, Incorporated, has been formed in Ithaca, New York. The founders of the company are Jeff Moskow, author of the popular, highly acclaimed Pas\(\mathrm{cal} / \mathbf{Z}^{(12)}\) compiler marketed by Ithaca InterSystems, Inc.; Laurie Hanselman Moskow, formerly Software Products Manager at InterSystems; and William Kellner, a software engineer who has worked extensively with Moskow on the Pascal/Z compiler.

The first product to be released by Software Building Blocks, Inc. is a two-pass, locally optimizing Pascal compiler for the IBM Personal Computer. The initial release will run under PC-DOS \({ }^{\text {®ve }}\); and a CP/M-86 \({ }^{(10)}\) version is planned for the near future. Based on the Pascal/ Z compiler, the Software Building Blocks implementation, SBB Pascal \({ }^{\text {(WW) }}\), closely follows the Jensen \& Wirth definition of the language, with extensions designed to aid the professional programmer in serious software development. Extensions will include: variable length strings, direct file access, arbitrary precision BCD numbers for business arithmetic, functions returning structured values, separate compilation, exterenal routines, include files, symbolic I/O of enumeration types, an ELSE clause for the CASE statement, overlays and chaining.

The compiler package includes a sophisticated interactive Pascal debugger, written in SBB Pascal, designed to aid in isolating and correcting faults in Pascal programs. Features of the debugger include the abilities to set and display both absolute and conditional breakpoints; set watches on variables, procedures or functions; display and modify both global and local variables; display the procedure/function stack, current statement and module numbers, current run-time requirements, and the last ten statements executed; trace through a program by statement number and procedure/function entry/exit; and more.

Also included in the package is a screen editor, provided in SBB Pascal source. The editor's capabilities include: insertion and deletion of lines and characters, finding and/or replacing of strings, copying lines of text, autoindent for entering structured programs, and many other features. The editor makes use of the function keys on the IBM PC to make editting as easy and efficient as possible. The editor is provided in source as an example of the advantages of programming in SBB Pascal. The library routines are also provided in 8086 assembly language source, and many other example programs are included as well.

Software Building Blocks, Inc. intends to release the PC-DOS version of the compiler in June. For more information, contact Laurie Moskow, Software Building Blocks, Inc., P.O. Box 119, Ithaca, New York, 14851-0119, (607) 272-2807.
(®10) \(\mathrm{CP} / \mathrm{M}-86\) is a trademark of Digital Research, Inc.
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\section*{SAGE OPENS BOSTON DIVISION}

Sage Computer Technology, headquartered in Reno, Nevada, has announced the opening of its Boston division.

The purpose of the new facility is to provide regional support for dealers and users of the Sage line of 16-bit microcomputers, and to expedite delivery of new units throughout the Eastern United States.

A complete inventory of Sage II's, Sage IV's, parts and literature is stocked, and a fully-equipped and staffed service department is maintained on the premises.

According to Rod Coleman, Sage president, plans call for a total of nine such offices to augment the company's domestic sales and support activities. 'Regional support for our dealers and OEMS is a critical part of our marketing plan.',

Sage's Boston office is now open to dealers \& OEMS, and is located at 15 New England Executive Park, Suite 120, Burlington, MA 01803. The telephone number is (617) 229-6868.

More information about Sage micros is available from either Boston office or corporate office at 4905 Energy Way, Reno, Nevada 89502. Telephone (702) 3226868.

If agency contact is required, phone or write The Schraff Group, 18226 W. McDurmott, Suite E, Irvine, CA 92714. Telephone (714) 540-8977.

\section*{NEW, 16-BIT SAGE IV HAS WINCHESTER PLUS MULTI-USER CAPABILITY}

RENO, NEVADA - Sage Computer Technology has announced availability of the Sage IV, 16-bit (68000) supermicro.

The new multi-user computer, which accommodates up to 6 simultaneous users, surpasses the considerable capabilities of the Sage II introduced in March, 1982.

Both machines are based on the 8 MHz 68000 processor, and both are capable of performing 2-million operations per second. According to Rod Coleman, Sage president, they offer performance comparable to that of high-end mini-computers at a mid-range to high-end business micro price.

The Sage IV comes standard with 128 K of main
memory which is expandable, optionally to a megabyte. This represents an enormous jump from the 128 K to 512 K expandability of the Sage II, which in turn offers far greater capacity than the typical \(64 \mathrm{~K}, 8\)-bit computer.

In addition, a 5 to 30 Mb Winchester disk, either fixed or removable, is built into the Sage IV next to a \(51 / 4\) inch floppy backup. Since there are no wait states, a 20 K program loads from the floppy in 1 second, and from the hard disk in \(1 / 10\) second.

The cabinet, though about \(11 / 2\) inches taller than that of the Sage II, is still deceptively small, measuring only \(61 / 2^{\prime \prime}\) high, \(12^{1} / 2^{\prime \prime}\) wide and \(163 / 4^{\prime \prime}\) deep.
"There aren't any tradeoffs with either of these machines, said Coleman, "the user doesn't have to give up software support to get high performance, because the Sages' p -System standard operating system is able to run hundreds of popular programs developed for 8bit micros."

More information may be had by contacting Sage Computer Technology, 35 North Edison Way, Suite 4, Reno, Nevada 89502. Telephone (702) 322-6868.

If agency contact is required, phone or write The Schraff Group, 1325 Airmotive Way, Suite 175, Reno, Nevada 89502. Telephone (702) 348-7339.

\section*{NEW MODULA-2 MANUAL FEATURES TUTORIALS, STANDARD LIBRARY}

DEL MAR, CA, Jan. 21 - A 264-page Modula-2 user's manual, featuring a language tutorial and standard library definitions, is now available from Volition Systems here.

Modula-2 is a new programming language designed by Niklaus Wirth to replace his earlier language, Pascal, in a wide range of real-world applications. Together with Wirth's own specifications of the language, this manual provides a complete description of Volition's implementation of Modula-2, according to its author Richard Gleaves of Volition Systems.

The manual is designed to be used with Wirth's 48page monograph which defines Modula-2 in a concise but informal style. The monograph is included with the manual. Wirth's newly published book Programming in Modula-2 is also available from Volition.

The manual contains a tutorial for Pascal programmers that can make them comfortable with the language within a few hours and proficient within a week, Gleaves said.

The name Modula-2 comes from MODUlar LAnguage. It uses modules to facilitate the development and maintenance of large, complex systems. The language is especially useful in large industrial and commercial applications where it can save software developers both time and money.

Modula-2 is designed to utilize standard software modules, which are defined in the new manual. These modules provide access to the facilities normally pro-
vided by an operating system, such as program and process control; console and file I/O, including random access files and disk directory operations; and storage management. The standard software modules also include utility routines for format conversion, strings, 19 digit BCD arithmetic, and other facilities.

The manual is divided into six sections. The Mod-ula-2 tutorial for Pascal programmers comprises about one-third of the book.

In addition, there is an introductory section and sections defining the standard library modules, the utility library, a system document that describes the implementation of Modula-2 for UCSD Pascal \({ }^{\text {(10) }}\) and a machine-specific implementation guide which includes information on machine specific library modules, interrupt handling, and machine-level data representation.

The Modula-2 User's Manual, including Wirth's Modula-2 report, is immediately available from Volition Systems, P.O. Box 1236, Del Mar, CA 92014 for \(\$ 35\) per copy. Wirth's book, Programming in Modula2, published in 1982 by Springer-Verlag, can be ordered for \(\$ 16\). Further information about the programming language is also available from Volition Systems.

Volition Systems concentrates on systems software development and on research and development in hardware and sotware. Since the company was founded in 1980, it has been a leader in the implementation and dissemination of the Modula-2 language and other high level languages and in the design and development of advanced computer architectures.

For further information, contact:
Volition Systems
P.O. Box 1236, Del Mar, CA 92014
(619) 481-2286
(al) UCSD Pascal is a trademark of the Regents of the University of California.

MODULA-2 USER'S MANUAL from Volition Systems (Del Mar, CA) describes Niklaus Wirth's new programming language in a 264 -page loose-leaf format. document contains a complete tutorial for Pascal programmers, sections defining the standard library modules and the utility library, and an implementation guide. The manual comes with a copy of Wirth's 48 -page technical report on Modula-2.

Modula-2 is particularly useful in large industrial and commercial applications where using standard modules facilitates development of large, complex systems, according to Volition, which has pioneered in commercial implementations of the new language. The Modula-2 User's Manual is immediately available from Volition Systems, P.O. Box 1236, Del Mar, CA 92014 for \(\$ 35\).

For further information, contact:
A. Winsor Brown
(714) 891-6043

\section*{USUS FALL MEETING SET FOR WASHINGTON,D.C.}

WASHINGTON, D.C., June 3 - USUS, Inc., the UCSD Pascal User's Society, will hold its semi-annual national meeting at the Crystal City Hyatt Hotel here October 14-16, according to Robert Peterson, USUS president.

In conjunction with the meeting, USUS will sponsor two free tutorials - an introduction to the p-System and an introduction to UCSD Pascal, including Apple Pascal.*

The meeting will feature technical presentations, hardware and software demonstrations, language tutorials, special interest group meetings and software library exchange. Also planned are expert user and major vendor panels. Election of officers will be held.
"Non-USUS members are welcome to register and attend any or all of the meeting programs,' Peterson noted.

USUS (pronounced use-us) represents users of the UCSD Pascal System and its derivatives including the UCSD p-System and Apple Pascal. It is the most widelyused, machine-independent software system. The society is non-profit and vendor independent.

The UCSD Pascal System has more than 100,000 users and is capable of running on nearly any computer. It was developed at the University of California San Diego to facilitate software portability.

Among the special interest group meetings scheduled for the Washington meeting are those for users of IBM Personal Computers, Apple, DEC, Texas Instruments, NEC Advanced Personal Computer, the IBM display writer and Sage Computer Technology computers.

Also meeting will be those interested in application development, graphics, communications, file access, Modula-2, UCSD Pascal compatibility and the Advanced System Editor.

The software library, with significant recent acquisitions, will be available for reproduction on various diskette formats. Members at the meeting will be able to copy the library onto their own disks for \(\$ 1.00\) each.

Those registering for the meeting before September 23 will qualify for the pre-registration price of \(\$ 25\). Checks should be made payable to USUS and mailed to USUS Meeting Committee, P.O. Box 1148, La Jolla, CA 92038. Registration at the door will be \(\$ 35\) and will begin at 10 a.m. Friday, October 14.

Hotel reservations should be made directly with the Crystal City Hyatt hotel (adjacent to Washington National Airport), 2799 Jefferson Davis Highway, Arlington, VA 22202, (703) 486-1234. Additional meeting information is available from Thomas Woteki, Information Systems Inc., 3865 Wilson Blvd., Suite 202, Arlington, VA 22203, (703) 522-8898.

USUS was created to promote and influence the development of the UCSD Pascal System and to provide users and vendors with a forum for education and information exchange about it. Annual membership in the society is \(\$ 25\) for individuals and \(\$ 500\) for institutions.

\footnotetext{
* Apple Pascal is a trademark of Apple Computer, Inc.
}

\section*{UCSD PASCAL USERS FORM TEXT EDITOR INTEREST GROUP}

SAN FRANCISCO, CA, June 15, 1983 - A special interest group (SIG) for users of the Advanced System Editor (ASE) for the UCSD Pascal System has been formed by USUS, the UCSD Pascal System User's Society, according to Robert W. Peterson, president of the society.

The new SIG will be chaired by Sam Bassett, of San Francisco, CA. "The ASE SIG will be open to any USUS member who is using or thinking about getting the Advanced System Editor." Peterson said.

The new ASE SIG allows members to share common problems and solutions and will serve as a clearing house for information relating to implementation, optimization and use of ASE on a variety of systems which have the UCSD p-System installed.

The SIG has established a liaison with Volition Systems of Del Mar, CA, the creators of ASE. It will coordinate relevant contributions to the USUS Software Exchange Library and to USUS News, the society's quarterly newsletter, Bassett said. Furthermore, SIG members may communicate via electronic mail under USUS sponsorship.

The next ASE SIG meeting will take place at the USUS semi-annual national meeting in Washington, D.C., October 14-16. In addition to the ASE and other SIG sessions, the USUS meeting will feature tutorials on UCSD Pascal and the UCSD p-System. Also on the agenda are technical presentations, software exchange, hardware and software demonstrations and an expert user panel.

Membership in the ASE SIG is free of charge to any member of USUS, the vendor-independent, nonprofit user's group for the UCSD Pascal System. Annual membership in the society is \(\$ 25\) for individuals and \(\$ 500\) for institutions.

USUS (pronounced use-us) was founded in 1980 to promote and influence the development of the UCSD Pascal System and to provide a forum for education and information exchange about it. Further information on USUS is available from the Secretary, USUS, P.O. Box 1148, La Jolla, CA 92038.

\section*{PASCAL USER'S SOCIETY FORMS MODULA-2 GROUP}

SAN DIEGO, CA, May 26 - USUS, the UCSD Pascal System User's Society, has formed a special interest group (SIG) for users of the new Modula-2 programming language, according to Robert W. Peterson, USUS president.

The new SIG will be chaired by David Ramsey of Corvus Systems, Inc. (San Jose, CA). The group was formed when USUS held its semi-annual national meeting here last month. Modula-2 runs on Version II based UCSD Pascal Systems.
"The Modula-2 SIG will be open to any USUS member using or wanting to investigate this language," Ramsey said. 'It is, to my knowledge, the first user's group devoted to communication about Modula-2."

The new language was created by Niklaus Wirth to answer difficulties encountered with his earlier language, Pascal. "As people discover the benefits of working in this new language, we expect this SIG to expand repidly,' Ramsey said.

Implementations of the Modula-2 programming language are available for the Apple II, //e and /// computers, the IBM Personal Computer, the 68000-based Sage 2 and 4, the Texas Instruments 9900, the Scenic One and Z80/8080-based systems, according to Joel J. McCormack of Volition Systems (Del Mar, CA).

Volition is the only current supplier of the language for use on microcomputers and supplies systems as well as the Modula-2 language to run on them.
"Because the language is modular, users spend less time writing and maintaining code," McCormack said. "Standard library modules provide Modula-2 with a standard operating environment, and programs created within it are portable across all Modula-2 systems."

The new Modula-2 SIG will enable users to share experiences with others using the language or developing applications in it, Ramsey said. "We expect to serve as a clearing house for user information in this fast-changing area."

One of the first goals of the SIG is creation of a user's library of Modula-2 programs that will be included in the USUS library, Ramsey noted. It will be compiled by Curt Snyder of Allergan Pharmaceuticals (Irvine, CA).

Membership in the Modula-2 SIG is free of charge to any member of USUS, which is the vendor-independent, non-profit user's group for the UCSD Pascal System. Annual membership in the society is \(\$ 25\) for individuals and \(\$ 500\) for institutions. Further information on USUS is available from the Secretary, USUS, P.O. Box 1148, La Jolla, CA 92038.

For those wanting to know more about the Mod-ula-2 SIG, Ramsey can be reached at Corvus Systems, 2029 O'Toole Avenue, San Jose, CA 95131, (408) 9467700, extension 267.

Volition Systems has pioneered in the implementation and dissemination of the Modula-2 language. Further information about Modula-2 and available implementations may be obtained from Tracy Barrett, Volition Systems, P.O. Box 1236, Del Mar, CA 92014, (619) 481-2286.

\section*{PASCAL USERS, VENDORS GATHER FOR USUS SAN DIEGO MEETING}

SAN DIEGO, CA, May 2 - USUS, the UCSD Pascal System User's Society, formed five new special interest groups (SIG's) and made plans for a first regional chapter at its well-attended, semi-annual national meeting here last week, according to Robert W. Peterson, USUS president.

In addition, two vendors of UCSD Pascal products - Apple Computer, Inc. and Volition Systems - chose the occasion to reveal new offerings.
"Record meeting attendance reflects the users' commitment to increased knowledge about use of the UCSD Pascal System,'" Peterson said. '"More than 240
attended and actively participated in special interest group and committee meetings, panel discussions and the four tutorials."

Keynote speaker for the event was Andrew Greenberg, designer and co-author of the popular Wizardry games. He told how he had solved the challenge of putting a very large program like Wizardry on a microcomputer with limited disk and main memory storage.
"Greenberg offered members valuable insights into program design, structure and implementation," Pe terson noted.

The new special interest groups are for application developers and for users of the NEC Advanced Personal Computer, the IBM Display Writer, the Advanced System Editor from Volition Systems, and the Modula-2 programming language. In addition, plans for the national organization's first local group in Southern California were discussed.

USUS already has SIG's for users of Apple, DEC, Texas Instruments and Sage computers, the IBM Personal Computer and for those interested in communications, word processing and UCSD Pascal compatibility.

Of particular interest to those attending the meeting was the demonstration area, where the latest advances in UCSD Pascal hardware, software and applications were demonstrated on 20 different machines, Peterson said.

At the meeting, Apple Computer, Inc., which has an installed base of some 82,000 Pascal development systems on its Apple II and Apple /// computers, announced that updates of Apple II Pascal and Apple /// Pascal will be available this year.

Apple revealed that Version 1.2 of Apple II Pascal will be available in the fourth quarter of 1983 and will provide support for all features of the Apple IIe including extended memory support for the 128 K IIe. Version 1.2 also makes available facilities for integrating into the UCSD Pascal environment in a natural way additional mass storage devices such as hard disks.

Apple also confirmed that Version 1.1 of Apple /// Pascal will be available at the end of June 1983. Its most notable feature is the Standard Apple Numeric Environment that fully implements the IEEE standard for floating point arithmetic.

Volition Systems demonstrated the new Modula-2 programming language running for the first time on an IBM Personal Computer. USUS members formed a Modula-2 SIG at the meeting to exchange information about the language. It will be chaired by Dave Ramsey, Corvus Systems (San Jose, CA).

The chairman of the newly formed application developer's SIG is Dennis Gallinat, Apple Computer (Cupertino, CA), and Samuel Bassett, Bassett Information Processing (San Francisco, CA) is chairing the Advanced System Editor SIG.

Lane Sharman, Resource Systems Group (Del Mar, CA) will head the Special Interest Group for the IBM Display Writer, and the NEC Advanced Personal Computer SIG will be chaired by George Symons, TICOM Systems, Inc. (Marina del Ray, CA).

The fall USUS meeting will be held in Washington, D.C., at the Hyatt Regency Crystal City, October 1416, 1983. Further information is available from the Sec-
retary, USUS, P.O. Box 1148, La Jolla, CA 92038.
USUS (pronounced use-us) is a vendor-independent, non-profit user's group for the most widely used, machine-independent software system - the UCSD Pascal System, and its successors such as the Apple Pascal System and the UCSD p-System.

USUS was created to promote and influence the development of the UCSD Pascal System and to provide a forum for education and information exchange about it. USUS has institutional as well as individual members in more than 20 countries. Annual membership in the society is \(\$ 25\) for individuals and \(\$ 500\) for institutions.

\section*{VOLITION DEMONSTRATES MODULA-2 FOR IBM PC}

DEL MAR, CA, May 3 - Volition Systems here has demonstrated Niklaus Wirth's new Modula-2 programming language running for the first time on the IBM Personal Computer.

The new implementation was demonstrated for members of USUS, the UCSD Pascal \({ }^{1}\) System User's Society, at its semi-annual national meeting in San Diego last week. Modula-2 will be included as part of Volition's complete software development system.
"Modula-2 is proving especially valuable in large industrial and commercial applications where standard software modules can save time and money in program development and maintenance,'" according to Joel J. McCormack of Volition Systems.
"Now our new implementation will make these savings possible on the IBM PC. Our software development system will even run efficiently on 64 K PC's," he continued. 'And the availability of Modula-2 on the IBM PC should make the language even more attractive to application developers."

The IBM PC implementation will significantly expand the availability of Modula-2. Current Volition versions are based on the 6502 (including Apple \(\mathrm{II}^{2}\) and Apple /// computers), the 8080/Z80, TI 9900, and the 68000.

Niklaus Wirth developed Modula-2 (from MODular LAnguage) to replace his earlier language, Pascal. Whereas Pascal was intended as a teaching language, Modula-2 is expressly designed for use in a wide range of real-world applications, and it offers great flexibility in the development of large, complex systems.

The implementation for the IBM PC is expected to be available in the third quarter of 1983, McCormack said. The system will include Modula-2 and Pascal compilers, the modula library, the powerful ASE text editor, V-NIX \({ }^{\text {© }}\) command shell (that provides a UNIX \({ }^{3}\) like programming environment), and a complete set of utility programs for file manipulation and electronic mail communication.

Volition Systems concentrates on systems software development and on research and development in hardware and software. Since the company was founded in 1980, it has led in the implementation and dissemination of the Modula-2 language and other high-level languages and in the design and development of advanced computer architectures.

For further information, contact:
Volition Systems
P.O. Box 1236, Del Mar, CA 92014
(619) 481-2286
\({ }^{1}\) UCSD Pascal is a trademark of the Regents of the University of California.
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\title{
Implementation Reports Implementation Reports Implementation Reports Implementation Reports \(\mathscr{I}_{\text {mplemen }}\)
}
0. DATE Apr. 28,1983
1. IMPLEMENTOR/MAINTAINER/DISTRIBUTOR /• Give a person, address and phone number •,
Robert Reimiller
OmegaSoft
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Camarillo, CA ..... 93010
(805) 987-6426
2. MACHINE/SYSTEM CONFIGURATION 1• Any known limits on the configuration or support software required, e. 9 MC 6809 Processor
Running Moos, 05-9, or Flex OS
Requires 48 K to 56 K (Recommended)
3. DISTRIBUTION /• Who to ask, how it comes, in what options. and at what price •)
North America: From Omega SoftInternational: From OmeqaSoft or distributors in Germany, Switzer-land, Great Iritain, Australia, Sweden, and the Netherlands. Priceis \(\$ 425\) to \(\$ 475\) for Compiler, Debugger, and Runtime.Relocatable Assembler/Linker available for \(\$ 125\) to \(\$ 150\).
4. DOCUMENTATION (*What is available and where •)
220 pg . Pascal manual with complete syntax and installationinstructions.
5. MAINTENANCE (• \(/ \mathrm{s}\) it unmaintained, fully maintained, etc) \({ }^{\bullet}\)
Yearly maintenance is \(\$ 100\) to \(\$ 125\)
6. STANDARD (• How does it measure up to standard Pascal? is it a subset) Extended) How •)Complete ISO standard except packed variables and procedural para-meters. Scored \(92 \%\) on conformance section of validation suite.ISO report in manual. Extended for real time and industrialcontrol applications.
7. MEASUREMENTS (• Of its speed or space •Marshals Algorithm: procedure size =270 bytes,lixecution time \(=9.7\) seconds
8. RELIABILITY (• Any information about field use or sites installed •, Over 400 sites installed. over 4000 sites installed.
9. DEVELOPMENT METHOD 1- How was it developed and what was it written in) •,
From scratch in assemble lanquare.
10 LIBRARY SUPPORT, (Any other support for compiler in the form of linkages to other languages source libraries. etc *Optional libraries to handle AMD9511 APU CHIP, and MuJti-TaskincPrimatives.

\section*{}

\section*{OmegaSoft Pascal Version 2}

\section*{Pascal Processor Identification}

Host Computer: Smoke Signal Broadcasting Chieftain 9522812 W 10 running the OS-9 operating system.

Host Computer Requirements: MC6809 processor, minimum of 48 K bytes of memory, 2 or more disk drives, running the OS-9, MDOS, XDOS, DOS69, or FLEX operating system.

Processor: OmegaSoft pascal version 2.10

\section*{Test Conditions}

Tester: R.D. Reimiller
Date: June 1982
Validation Suite Version: 3.0
General Introduction to the OmegaSoft Implementation
The OmegaSoft Pascal compiler was developed to provide the users of the 6809 processor with a fast and efficient way to develop code capable of running on the host development system or installed into a target system. The compiler is aimed primarily at industrial applications such as process control and instrumentation. Due to the nature of these applications many extensions were added such as byte arithmetic, long integers, dynamic length strings, modular compilation, and versatile variable addressing. As a secondary requirement it was desired that the compiler be able to accept a Pascal program written in ISO standard Pascal wherever possible.

\section*{CONFORMANCE TESTS}

Number of tests passed \(=144\)
Number of tests failed \(=12\) ( 9 reasons)

\section*{Details of Failed Tests}

Test 6.4.2.3-3: If an enumerated type is defined in the index declaration part of an array its values cannot be referenced until the array declaration is complete.

Test 6.4.2.3-4: If an enumerated type is defined in a record its values cannot be referenced until the record declaration is complete.

Tests 6.6.3.1-4, 6.6.3.4-1, 6.6.3.4-2, and 6.6.3.5-1: Procedures and functions cannot be passed as parameters.

Test 6.6.5.4-1: Pack and Unpack procedures are not supported.

Test 6.7.2.2-3: Failed on MOD using a negative dividend. The Jenson/Wirth 'remainder after division'' method is used rather than the method specified in the ISO standard.

Test 6.8.2.4-1: Non-local GOTO's are not allowed.
Test 6.8.3.9-1: Assignment to the control variable of a FOR loop occurs after the evaluation of the first expression.

Test 6.9.3-1: Standard I/O devices may not be redefined if declared.

Test 6.9.3.5.1-1: Real numbers written out in floating point format always have six digits to the right of the decimal point.

\section*{DEVIANCE TESTS}

Number of deviations correctly detected \(=83\)
Number of tests showing true extensions \(=45\) (22 reasons)
Number of tests not detecting erroneous deviations \(=\) 9 (6 reasons)

\section*{Details of Extensions}

Test 6.1.5-4: No digits are needed after the decimal point in a real number.

Tests 6.1.6-4 and 6.1.6-5: Labels may be a positive integer constant.

Tests 6.1.7-5, 6.4.3.1-3, 6.4.3.1-4, 6.6.3.3-5, 6.9.3.22: All variables are packed at the byte level, the reserved word 'Packed'" is ignored in any type declaration.

Tests 6.1.7-6, 6.1.7-7, 6.1.7-8, 6.4.3.2-5: Strings, characters, and arrays of less than 127 elements are all compatible.

Tests 6.1.7-11 and 6.4.5-12: Strings are dynamic length, allowable length is from 0 (null string) to 126.

Tests 6.2.1-8 and 6.2.1-10: Label, cont, type, and var declaration sections can be in any order and reprated multiple times until a procedure/function declaration or "begin'" is encountered.

Test 6.3-9: In any context where a constant is acceptable an expression with a constant value may be used.

Test 6.4.2.3-5: All enumerated type values are compatible.

Test 6.4.3.3-8: The values of the case constants in a record variant declaration are not used, access is provide to all variants at all times.

Test 6.4.5-7: All subranges of the same type are compatible.

Tests 6.4.5-8 and 6.4.5-13: Arrays of the same size are compatible.

Tests 6.4.5-9 and 6.4.6-7: Records of the same size are compatible.

Test 6.4.5-10: All pointers are compatible with other pointers or the type "Hex'".

Test 6.6.2-5: Any type with a size of less than 128 bytes can be used as a function return type.

Test 6.6.6.3-2: Trunc and round can have integer or longinteger parameters.

Test 6.7.2.3-2: Logical operators are valid for charaster and integer expressions.

Test 6.7.2.5-6: Arrays of the same size can be compared. Records of the same size can be compared.

Test 6.8.2.4-2: Goto between branches of an If statement are allowed.

Test 6.8.2.4-3: Goto between branches of a Case statement are allowed.

Tests 6.8.3.5-7 and 6.8.3.5-8: Subrange Case statement constants are allowed.

Tests 6.8.3.9-5, 6.8.3.9-6, 6.8.3.9-7, 6.8.3.9-10, 6.8.3.9-12, 6.8.3.9-13, 6.8.3.9-14, 6.8.3.9-15, 6.8.3.9-16, and 6.8.3.9-17: No restrictions are placed on For statement control variable.

Tests 6.8.3.9-8 and 6.8.3.9-9: If a For statement is entered and exited normally the control variable will be valid and contain the final value. If a For statement is not entered then the control variable will be valid and contain the initial value.

\section*{Details of Deviations}

Test 6.1.8-5: A number can be terminated by a letter.
Tests 6.2.1-5 and 6.2.1-6: Multiple siting for labels is not checked, nor are labels required to be sited at all.

Tests 6.2.2-8, 6.3-6, and 6.4.1-3: Error in scope rules.

Test 6.6.1-7: Unresolved forward function or procedure declaration is not detected.

Test 6.6.3.3-4: Use of a field selector as a parameter is not detected.

Test 6.10-4: No check is made for duplication of program parameters.

\section*{ERROR-HANDLING}

Number of errors correctly detected \(=19\)
Number of errors not detected \(=31\) (13 reasons)

\section*{Details of Errors Not Detected}

Tests 6.2.1-11, 6.4.3.3-11, 6.4.3.3-12, 6.4.3.3-11, 6.5.4-2, and 6.6.2-9: No checking is made to verify whether or not a variable is accessed that has an undefined value. Instead the variables are guaranteed to contain garbage unless initialized.

Tests 6.4.3.3-1, 6.6.5.3-8, 6.6.5.3-9, and 6.6.5.3-10: Any tagfields or selector variables in a record variant are irrelevant to which variants can be accessed.

Test 6.4.6-10: No subrange checking on parameter passing.

Tests 6.4.6-12, 6.4.6-13, and 6.7.2.4-4: Overflow checking is done on sets based on byte count - not per element.

Tests 6.5.4-1, 6.6.5.3-4, 6.6.5.3-5, and 6.6.5.3-11: Pointer value is not checked before use.

Tests 6.5.5-2, 6.5.5-3, 6.6.5.3-6, and 6.6.5.3-7: There are no restrictions on the use of pointers or file buffer variables which are currently parameters or elements of a with statement.

Test 6.6.5.2-5: To support random files a "get'" is not executed until called as a procedure or when accessing the file buffer without a valid element - not at the time of "reset".

Test 6.6.6.4-7: Char and Hex variables "roll over" from maximum value to zero - it is not considered an error.

Test 6.6.6.5-7: If eof is true - so is eoln - it is not
considered an error to check eoln if eof is true.
Tests 6.8.3.5-10 and 6.8.3.5-11: If no match in case statement, falls through with no error.

Test 6.8.3.9-18: No restrictions on the control variable of a For loop.

Test 6.8.3.9-1: At the completion of a For loop the control variable is valid and has the final value.

Tests 6.9.3.2-5 and 6.9.3.2-5: Writing of real numbers with no digits past the decimal point is permissible.

\section*{QUALITY MEASUREMENT}

Number of tests run \(=52\)
Number of tests incorrectly handled \(=5\)

\section*{Results of Tests}
"Synthetic Benchmark" - execution time 1 minute, 10 seconds.
"GAMM measure" - execution time 1 minute, 40 seconds for \(\mathrm{N}=1000\)
procedure calls - execution time 40 seconds
identifiers are significant up to 120 characters.
source lines may be up to 120 characters.
no reasonable limit on number of real literals allowed.
no reasonable limit on number of strings allowed.
if a line of code is incorrectly part of an unclosed comment the compiler will signal that no code was generated for the line.
at least 50 types may be declared in a program.
no reasonable limit on number of labels, but there can be a maximum of 8 forward referenced goto's in a block.
at least 128 constant definitions are allowed per constant declaration part.
at least 128 procedures are permitted in a program.
maximum size for an array or record or for any variable section is 32750 bytes.
at least 8 index types can appear in an array type.
at least 128 case-constant values are permitted in a variant record.
at least 50 record-sections can appear in the fixed part of a record.
at least 30 distinct variants are permitted in a record.
''Warshall's algorithm'' procedure size \(=270\) bytes, execution time \(=9.7\) seconds.
considerably less than 300 indentifiers are allowed in a declaration list (actual number depends on length of identifier).
at least 8 dimensional array is allowed.
procedures may be nested to at least 15 levels.
at least 30 formal parameter sections can appear in one parameter list.
the dispose in the standard heap manager is a dummy, a more complex heap manager is available.
deeply nested function calls are allowed (at least \(6)\).
deeply nested compound statements are allowed (at least 25).
a procedure may have at least 300 statements.
deeply nested if statements are allowed (at least 25).
at least 256 case constants are allowed.
at least 300 constants are allowed in a case-constant list.
case statements can be nested to at least 15 deep. repeat loops can be nested to at least 15 deep. while loops can be nested to at least 15 deep. for loops can be nested to at least 15 deep.
with statements can be nested to at least 15 deep. recursive I/O can be used with the same file for the second I/O action.
at least 30 variable-accesses can appear in a read or readln parameter list.
at least 30 write-parameters can appear in a write or writeln parameter list.
data written on the output field appears regardless of the omission of a line marker.

\section*{IMPLEMENTATION-DEFINED}

Number of tests run \(=12\)
Number of tests incorrectly handled \(=1\)

\section*{Details of Implementation-Defined Features}

Tests 6.1.9-5 and 6.1.9-6: alternate symbols are available for comments, array indices, and pointers.

Test 6.4.2.2-10: Maxint is 32767
Test 6.4.3.4-5: maximum range of set elements is \(0 . .1007\)

Test 6.6.6.2-11: Base \(=2\), Bits of mantissa \(=24\), not rounding, minimum value \(=2.710506 \mathrm{E}-20\), maximum value \(=9.223372 \mathrm{E}+18\)

Tests 6.7.2 3-3 and 6.7.2.3-4: Boolean expressions
are fully evaluated.
Tests 6.8.2.2-1 and 6.8.2.2-2: In an assignment statement evaluation of the expression is done before the selection of the variable.

Test 6.8.2.3-2: When a procedure is called the parameters are evaluated in forward order.

Test 6.9.3.2-6: Default field widths are: Integers \(=\) 10, Boolean \(=6\), Real \(=16\), Longinteger \(=16, \mathrm{Hex}=\) 6.

Test 6.9.3.5.1-2: Real values written in floating point format have 2 exponent digits.

Test 6.9.3.6-1: Boolean values written in the default fieldwidth have the format as shown (between quotes) " TRUE" and " FALSE".

\section*{Details of Tests Incorrectly Handled}

Tests 6.6.6.1-1: Functions are not allowed to be passed as parameters to a procedure.

Level 1 Tests - Not applicable

\section*{EXTENSIONS}

Extension present \(=1\)

\section*{Result of Extension}

Test 6.8.3.5-16: An otherwise clause is allowed on a case statement.

\section*{Pascal News}

2903 Huntington Road Cleveland, Ohio 44120

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\section*{1 Magnetic tape}
\(9-T r a c k\), odd parity. \(1 / 2^{\prime \prime} \times 600\) '. Select Density:
( ) 800 bpi ( ) 1600 bpi
1 ) ANSI-STANDARD. Each logical record is an 80 character card image. Each physical record has a block size of \(\mathbf{4 0}\) logical records. Select Character Code:
( ) ASCII | ) EBCDIC
( ) Special DEC System Alternate Formats: ( ) RSX-IAS PIP (requires ANSI MAGtape RSX SYSGEN). 1 ) DOS-RSTS FLX.

\section*{\(11^{\prime \prime}\) Diskette}
( ) Single Density
( ) Double Density

\section*{Format}
( ) CP/M I J UCSD III (W. D. Microengine)
1 I UCSD II. IV I I DEC-RT (Single Density)
( J DEC-RSX Files 11 | | IBM 3740 (Single Density EBCDIC)

\section*{Special Format}
( ) Interieave (1-26)
1 ) Skew (0-25)

Office Use Only

Signed:
Date:

\section*{Richard J. Cichelli}

On Behalf of A. H. J. Sale and B. S. I.

\section*{Us}

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\section*{Facts about Pascal, THE PROGRAMMING LANGUAGE:}

Pascal is a small, practical, and general-purpose (but not all-purpose) programming language possessing algorithmic and data structures to aid systematic programming. Pascal was intended to be easy to learn and read by humans, and efficient to translate by computers.

Pascal has met these goals and is being used successfully for:
- teaching programming concepts
- developing reliable "production" software
- implementing software efficiently on today's machines
- writing portable software

Pascal implementations exist for more than 105 different computer systems, and this number increases every month. The "Implementation Notes" section of Pascal News describes how to obtain them.

The standard reference ISO 7185 tutorial manual for Pascal is:
Pascal - User Manual and Report (Second, study edition)
by Kathleen Jensen and Niklaus Wirth.
Springer-Verlag Publishers: New York, Heidelberg, Berlin
1978 (corrected printing), 167 pages, paperback, \$7.90.
Introductory textbooks about Pascal are described in the "Here and There" section of Pascal News.
The programming language, Pascal, was named after the mathematician and religious fanatic Blaise Pascal (1623-1662). Pascal is not an acronym.

Remember, Pascal User's Group is each individual member's group. We currently have more than 3500 active members in more than 41 countries.

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