# Pascal News 

December, 1978 NUMBER 13 (oh, how unlucky...)
communications about the programming language Pascal by Pascalers

Cover
0
1
3
5
5
6
8
9
10
12
16
18
31
32
34
59
59

* Pascal News is the official but informal publication of the User's Group.

Pascal News contains all we (the editors) know about Pascal; we use it as the vehicle to answer all inquiries because our physical energy and resources for answering individual requests are finite. As PUG grows, we unfortunately succumb to the reality of (1) having to insist that people who need to know "about Pascal" join PUG and read Pascal News - that is why we spend time to produce it! and (2) refusing to return phone calls or answer letters full of questions - we will pass the questions on to the readership of Pascal News. Please understand what the collective effect of individual inquiries has at the "concentrators" (our phones and mailboxes). We are trying honestly to say: "we cannot promise more than we can do."

* An attempt is made to produce Pascal News 3 or 4 times during an academic year from July 1 to June 30; usually September, November, February, and May.
* ALL THE NEWS THAT FITS, WE PRINT. Please send material (brevity is a virtue) for Pascal News single-spaced and camera-ready (use dark ribbon and 18.5 cm 1 ines!).
* Remember: ALL LETTERS TO US WILL BE PRINTED UNLESS THEY CONTAIN A REQUEST TO THE CONTRARY.
* Pascal News is divided into flexible sections:

POLICY - tries to explain the way we do things (ALL PURPOSE COUPON, etc.). EDITOR'S CONTRIBUTION - passes along the opinion and point of view of the editor together with changes in the mechanics of PUG operation, etc.
HERE AND THERE WITH PASCAL - presents news from people, conference announcements and reports, new books and articles (including reviews), notices of Pascal in the news, history, membership rosters, etc.
APPLICATIONS - presents and documents source programs written in Pascal for various algorithms, and software tools for a Pascal environment; news of significant applications programs. Also critiques regarding program/algorithm certification, performance, standards conformance, style, output convenience, and general design.
ARTICLES - contains formal, submitted contributions (such as Pascal philosophy, use of Pascal as a teaching tool, use of Pascal at different computer installations, how to promote Pascal, etc.)
OPEN FORUM FOR MEMBERS - contains short, informal correspondence among members which is of interest to the readership of Pascal News.
IMPLEMENTATION NOTES - reports news of Pascal implementations: contacts for maintainers, implementors, distributors, and documentors of various implementations as well as where to send bug reports. Qualitative and quantitative descriptions and comparisons of various implementations are publicized. Sections contain information about Portable Pascals, Pascal Variants, Feature-Implementation Notes, and Machine-Dependent Implementations.

* Volunteer editors are (addresses in the respective sections of Pascal News):

Andy Mickel - editor
Jim Miner, Tim Bonham, and Scott Jameson - Implementation Notes editors
Sara Graffunder and Tim Hoffmann - Here and There editors
Rich Stevens - Books and Articles editor
Rich Cichelli - Applications editor
Tony Addyman and Rick Shaw - Standards editors
Scott Bertilson, John Easton, Steve Reisman, and Kay Holleman - Tasks editors

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+ reproduce, etc and
$\leftarrow$.
* mail to this address.
/ / Please enter me as a new member of the PASCAL USER'S GROUP for $\qquad$ Academic year(s) ending June 30, $\qquad$ (not past 1982). I shall receive all the issues of Pascal News for each year. Enclosed please find $\qquad$ . (* Please see the POLICY section on the reverse side for prices and if you are joining from overseas, check for a PUG "regional representative." *)
/ / Please renew my membership in PASCAL USER'S GROUP for $\qquad$ Academic year(s) ending June 30, $\qquad$ (not past 1982). Enclosed please find $\qquad$ .
/ / Please send a copy of Pascal News Number(s) $\qquad$ . (* See the Pascal News POLICY section on the reverse side for prices and issues available. *)
/ / My new $\begin{gathered}\text { address } \\ \text { phone }\end{gathered}$ is printed below. Please use it from now on. I'll enclose an old mailing label if I can find one.
(* The U.S. Postal Service does not
/ / You messed up my address. See below. forward Pascal News. *)
/ / Enclosed please find a contribution (such as what we are doing with Pascal at our computer installation), idea, article, or opinion which I wish to submit for publication in the next issue of Pascal News. (* Please send bug reports to the maintainer of the appropriate implementation listed in the Pascal News IMPLEMENTATION NOTES section. *)
/ / None of the above. $\qquad$
$\qquad$
$\qquad$
$\qquad$
Other comments:
From: name $\qquad$
mailing address $\qquad$
$\qquad$
$\qquad$
$\qquad$
phone $\qquad$
computer system(s) $\qquad$
date $\qquad$
(* Your phone number aids communication with other PUG members. *)

JOINING PASCAL USER'S GROUP?

- membership is open to anyone: particularly the Pascal user, teacher, maintainer, implementor, distributor, or just plain fan.
- please enclose the proper prepayment (checks payable to "Pascal User's Group"); we will not bill you.
- please do not send us purchase orders; we cannot endure the paper work! (If you are trying to get your organization to pay for your membership, think of the cost of paperwork involved fọr such a small sum as a PUG membership!)
- when you join PUG anytime within an academic year: July 1 to June 30 , you will receive all issues of Pascal News for that year unless you request otherwise.
- please remember that PUG is run by volunteers who don't consider themselves in the "publishing business." We produce Pascal News as a means toward the end of promoting Pascal and communicating news of events surrounding Pascal to persons interested in Pascal. We are simply interested in the news ourselves and prefer to share it through Pascal News, rather than having to answer individually every letter and phone call. We desire to minimize paperwork, because we have other work to do.
- American Region (North and South America): Join through PUG(USA). Send $\$ 6.00$ per year to the address on the reverse side. International telephone: 1-612-376-7290.
- European Region (Europe, North Africa, Western and Central Asia): Join through PUG(UK). Send $£ 4.00$ per year to: Pascal Users' Group/ c/o Computer Studies Group/ Mathematics Department/ The University/ Southampton S09 5NH/ United Kingdom. International telephone: 44-703-559122 x700.
- Australasian Region (Australia, East Asia - incl. Japan): Join through PUG(AUS). Send $\$$ A8.00 per year to: Pascal Users Group/ c/o Arthur Sale/ Dept. of Information Science/ University of Tasmania/ Box 252C GP0/ Hobart, Tasmania 7001/ Australia. International Telephone: 61-02-23 0561.
PUG(USA) produces Pascal News and keeps all mailing addresses on a common list. Regional representatives collect memberships from their regions as a service, and they reprint and distribute Pascal News using a proof copy and mailing labels sent from PUG(USA). Persons in the Australasian and European Regions must join through their regional representatives. People in other places can join through PUG(USA).
RENEWING? (Costs the same as joining.)
- please renew early (before August) and please write us a line or two to tell us what you are doing with Pascal, and tell us what you think of PUG and Pascal News to help keep us honest. Renewing for more than one year saves us time.
ORDERING BACKISSUES OR EXTRA ISSUES?
- our unusual policy of automatically sending all issues of Pascal News to anyone who joins within an academic year (July 1 to June 30) means that we eliminate many requests for backissues ahead of time, and we don't have to reprint important information in every issue--especially about Pascal implementations!
- Issues 1, 2, 3, and 4 (January, 1974 - August, 1976) are out of print.
- Issues 5, 6, 7, and 8 (September, 1976 - May, 1977) are out of print. (A few copies of issue 8 remain at PUG(UK) available for $\hbar 2$ each.)
- Issues 9, 10, 11, and 12 (September, 1977 - June, 1978) are available from PUG(USA) all for $\$ 10$ and from PUG(AUS) all for \$A10.
- extra single copies of new issues (current academic year) are: $\$ 3$ each - PUG(USA); $£ 2$ each - PUG(UK); and \$A3 each - PUG(AUS).

SENDING MATERIAL FOR PUBLICATION?

- check the addresses for specific editors in Pascal News. Your experiences with Pascal (teaching and otherwise), ideas, letters, opinions, notices, news, articles, conference announcements, reports, implementation information, applications, et.c. are welcome. "All The News That Fits, We Print." Please send material single-spaced and in camera-ready (use a dark ribbon and lines 18.5 cm wide) form.
- remember: All letters to us will be printed unless they contain a request to the contrary.


## MISCELLANEOUS INQUIRIES?

- Please remember that we will use Pascal News as the medium to answer all inquiries, and we regret to be unable to answer individual requests.

Ln
$\begin{array}{ll}\text { UNIVERSITY OF MINNESOTA } & \begin{array}{l}\text { University Computer Center } \\ \text { TWIN CITIES }\end{array} \\ 227 \text { Experimental Engineerin }\end{array}$ University Computer Center
227 Experimental Engineering Building (612) 376-7290
(* This is going to be a long column. I apologize, but many important things need to be said. The future of PUG is one of them! *)

I would like to thank everyone who has helped with Pascal User's Group and Pascal News. Three far-sighted individuals to whom we owe special thanks are close by at the University Minnesota). Pete Patton is our Computer Center director and Larry Liddiard is our ssociate director for systems. Phil Voxland is the director of the Social Science Research Facilities Center. Their encouragement and moral support gave us the chance to see Pascal through to its widespread acceptance through the medium of Pascal News.

## 0. FORTRAN

Being a member (just joking) of ACS and SHAFT (American COBOL Society--dedicated to the elimination of COBOL in our lifetime--and the Society to Help Abolish FORTRAN Teaching, like to add that if the new FORTRAN compilers are written in assembler like most of the old nes, then we should see instability as well. Pascal may make its move on the large machine especially at Universities!

FORTRAN - The End at Last?

## D.W. Barron

The apparent indestructibility of FORTRAN as the preferred programming language of users in the physical sciences has long been a discouragement to those of us who try to spread the use of Pascal. We have thought long and hard about ways to convert the FORTRAN f Pascal News probably don't follow the activities of the FORTRAN Standards Committee, and so will be unaware of recent developments which indicate that the Standards Committee is doing the job for us. The specification of FORTRAN 77 has recently been published -
a hotpotch of "features" heaped indiscriminately on the existing FORTRAN language in a way that is not downwards compatible. The Committee has already started work on FORTRAN 82 and has published a
draft list of features of the "central module". Since such sacred things as $C$ in column 1 for comments, continuation in column 6 and tatements starting in column 7 are apparently to go, the result will not even bear a superficial resemblance to the FORTRAN that present day users know and love

If you ask a scientist why he uses FORTRAN, his answer will include some or all of the following reasons
i) it is efficient
ii) it is simple
iii) it is universal.

Editor's Contribution

The first reason is a red-herring - Pascal is probably more efficien than FORTRAN on many computers, but the typical FORTRAN user has been brainwashed into believing that nothing can be more efficient than FORTRAN. The simplicity of FORTRAN is superficial - true simplicity comes from logical cohesion. Certainly, FORTRAN 77 can't be describe card in the FORTRAN pack is its universality. Every computer centre has a FORTRAN compiler, they are reasonably compatible, and the scientist can move his FORTRAN programs from place to place with relative ease. The reason for this is that FORIRAN has been around for a long time, and has been stable for a long time. It is this stability that the new Standards are destroying. Now, FORTRAN ma be ANSI or ' 77 . ' 82 lurks in the wings. These new versions are appreciably different from the old, so FORTRAN loses its identity.
What gives a language an identity? Whability. The FORTRAN philosophy seems to be that FORTRAN is anything to which a particular committee chooses to give that name at a stroke they have destroyed FORTRAN's most valuable assett.

Whilst rejoicing over this development, we should not lose sight of the moral for Pascal. When it comes to determining the usage of a language, having a good language helps, but most important is to have a stable, widely available language. That way we can reach hery language then we must deny ourselves the indulgence of changing
I. Recent Events (at least since PUGN \#12)

A lot of people responded positively to the new Applications section started in \#12. I ope this issue's Applications section is just as worthy. There were also a few comments in favor of regularly featuring "Pascal in Teaching."

It's been quite a while since \#12 appeared and even though we have been flooded with renewals and the enthusiastic remarks of "keep-up-the-good-work!", here I am putting \#13 out very ate. I'm sorry. "13" is turning out to be unlucky indeed. Please see part II

Part of the reason we're late is that it is hard to keep up with the swirl of events
surrounding Pascal. So... case ImportantEvent of

## Employment:

Please see the letter from Chuck Beauregard in the Open Forum section. People have been
calling me constantly on the phone for 6 months now trying to find Pascal people to fill
jobs they offer. So get the word out: IF YOU KNOW Pascal YOU
ConcurrentPascal:
Per Brinch Hansen is trying to survey Concurrent Pascal users. Please respond to his
Fetter:†日 Ehe Implementation Notes section before February 28
NASA:
The United States National Aeronautics and Space Administration (NASA) is making a strong committment to Pascal. NASA Langley, NASA Ames, and NASA in Houston are all concerned with reltable software in deep-space probes (such as the upcoming Galileo project).
member John Knight (who is the CDC-Star Pascal implementor and convenor of the join member John Knight (who is the CDC-Star pascal implementor and convenor of the joint
SIGPLAN/PUG session at ACM--see below) has been keeping us informed. Thanks, John!
ConventionalizedExtensions:
In \#12 we described the formation of an International Working Group on Pascal Extensions which is supposed to decide on a handful of conventionalized extensions. Please see the section in Open Forum on Standards.

## Editor's Contribution

Standards:
The British Standards Institute Working Group (DPS/13/4) work on a Pascal standard (see PUGN \#14 for a working draft) should help lay to rest much of the standards controversy. A revised version of the document will be accepted by ISO (and therefore ANSI in the United States) in mid-1979. Politics with standards unfortunately keeps growing as knowledge and use of Pascal grows. And unfortunately ANSI has decided to refer

## PascalMachines:

One of our fondest wishes has been that hardware manufacturers help bridge the gap to Pascal by building machines with friendlier architectures. In fact, long-time PUG member Judy Mullins Bishop wrote her PhD thesis under Prof. David Barron investigating just such a Pascal architecture, which among other things, would require minimal storage
requirements for object code. Some people keep saying that BASIC (BASICK) exists on small personal computers and that Pascal implementations are too big. The fact is that the popular smal 1 personal computers are now based on microprocessors such and to do much useful anyway!). Thus a Pascal implementation is at a disadvantage having to emulate actions that should be performed in the hardware to begin with, and therefore consuming more code space. The significant aspect of the widely-known UCSD (University of Calif. at San Diego) Pascal project was to dispel the myth that Pascal couldn't run on a micro. However, in order to be small, this implementation had to be kept interpretive (there exist several cross-compilers of "hard code" for these micros from other sources). Also because UCSD Pascal is a Pascal-P derivative, the P-code had to be modified and packed (ffrequency-encoded). The result is slower execution.
But, recently, Western Digital built an inexpensive chip-set expressly for running the modified UCSD P-code, and a speedup of 5 or 10 is being realized. National Semiconductor wtandard P-code that will even better in a few months by building something closer to wraps for over two years! After being disappointed by Zilog and the Z-8000 a year ago, it is good to see Western Digital take the courageous first step, made possible by the people at UCSD. (Unfortunately we have been receiving altogether too many reports from users that UCSD Pascal is not as stable as it should be, and that its non-standard extensions are particularly lacking in robustness. For an example, see the Implementation Notes section.)

## PascalUsage:

The Western Digital product brochure for its Pascal "Microengine" apparently misquoted Ken Bowles of UCSD concerning: "there are more users of UCSD's Pascal today than users of all other versions combined." This patently false statement has caused Ken some 8\% (1 assment, and although it's hard to get an exact figure, I d estimate that near UCSD) followed by IBM 370 's followed by CDC machines and DEC 10 's and 20 's. To give specific example, the venerable CDC-6000 implementation is running at over 300 (very large) sites, and at just one of them (our University of Minnesota computer center) the compiler was accessed over 272,000 times from 77/07/01 to 78/06/30 which represented a $68 \%$ increase over the previous 12 months. We have been trying to collect usage data through the checklists in the Implementation Notes section and will try to summarize them in one place in a future issue.

ExplosionInIndustryLiterature:
Byte, Electronics, Creative Computing (ROM), and others have run full-1ength articles on Pascal. In fact the August, 1978 Byte was almost entirely devoted to Pascal! This phenomenon is most encouraging because eventually the mainstream computer literature languages. Other computer journals (Computer World, for example) have kept Pascal in the news this last six months and we appreciate it. The only bad side effect is that the publicity has literally swamped us here at PUG central with mail and phone calls.

## I. Pascal User's Group / Pascal News status

## Running Out of Time

Just at the time when the mail was starting to build up last May, (it now runs between 10 and 30 pieces a day), our usually smooth-running operation became short-handed. Jim Miner started going to school full-time. Sara Graffunder delivered a 2.81 kg baby boy named David. (As an aside, Rich Stevens got married two days after Thanksgiving in and UCSD workshop politics all began to consume our time with very little in return (just working very hard to stay still). The cover of this issue depicts the situation.

As if that weren't all, I do have my own full-time job to do here at the University of Minnesota computer center. This summer we changed operating systems and character sets Because I am also involved with the project to produce a new release of CDC-6000 Pascal was unable to work on PUG much at all this summer. Our mail went unanswered, and I apologize.

## New Members

This is the first academic year (July, 1978 to June, 1979) for almost 1000 new members, and I wanted to assure them good service and information which would tell them what our style is like. But after catching up with the mail in October (and returning all $\$ 4$ renewals arriving after August 1 thus allowing a 1 -month grace period) and falling behind again, think I have disappointed quite a few people and I'm sorry. We have stated that we are all-volunteer, and that we have little or no secretarial help, but you new members have yet to read this sentence because you have received nothing from us until now! I urge new members to get backissues from last year--see the section on backissues in Here and There.

## Deadlines

We have received some sharp criticism from overseas PUG members (who, by the way contribute most of the material for publication!) about the publications deadlines for Pascal News. The fact is that we had no deadlines during 76-77 (issues 5-8) and everything
 an issue, overseas members often received issues after the deadline for the next in producing Solution: lets go Solution: let's go back to no deadlines. If you have material, simply send it in. Confusion

Our mailing list has never been sold or given out. Any PUG members with issues of Pascal News from \#9/10 onward has the mailing list, because we print the roster. We have however sent out a notice last month for the jointly-sponsored ACM SIGPLAN (Association for Computing Machinery Special Interest Group on Programming Languages)-PUG session at the national ACM 78 conference this December, and it is already causing confusion. We didn received from us. If I were on the receiving end I would be confused too! We knew we were going to be late with this issue, and that is why we sent the notice out.

## Summary

I hate to paraphrase someone like Winston Churchill, but he said that sometimes doing your best is not enough--sometimes you have to do what is required. Please read on in my open letter in the Open Forum section.

TIDBITS
Ole Anderson, Corvallis, OR 97330: "I have a LISP interpreter that runs under the UCSD Pascal system- Would anyone be interested?" (*78/05/19*)

David A. Beers, Santa Ana, CA 92701: "I very much enjoy reading Pascal News. It is a refreshing exposure to rationality when compared to my job as a business systems programmer. ... I have talked to Joseph Mezzaroba of Villanova Universiy concerning
his DOS/VS version of the AAEC Pascal 8000, and will be attempting to convert it to DOS his DOS/VS version of the AAEC Pascal 8000, and will be attempting to convert it to DO unless I hear of someone else's successful endeavors in this area." ( $* 78 / 10 / 25 *$ )
C. Y. Begandy, Aluminum Company of America, Alcoa Center, PA 15069: "I recently obtained the Pascal compiler from the DECUS library. Because of daytime core usage restrictions
at our installation, it is necessary to decrease the size of the executable program. at our installation, it is necessary to decrease the size of the executable program
Any information you might have on other users ${ }^{\prime}$ experiences in implementing either a Any information you might have on other users experiences in implementing either a
smaller version of this compiler, or a segmented version would be greatly appreciated." (*78/05/26*)
Gerd Blanke, Eschborn, Germany: "... MODULA will be running on a zilog MCS with 64 K under Rio near the end of this year!" ( $* 78 / 10 / 27 *$ )

John H. Bolstad, Department of Mathematics, Florida State Univ, Tallaḥassee, FL 32306: "We use Pascal here for almost all computer science courses. The system programmers also use it." (*78/07/11*)
R. T. Boute, Francis Wellesplein 1, B-2000 Antwerpen Belgie: "We are interested in a special hardware support for standard and concurrent Pascal, for example microprogrammed implementations of the P machine." ( $* 78 / 10 / 17 *$ )
Robert Boylan, Metromation, Princeton, NJ 08540: "I know a PDP-11 version of Pascal is in existence, but has anyone done one for a Modcomp mini?" (*78/07/26*)

David C. Cline, Westboro, Mass 01581: "Pascal is attracting a lot of attention here at Data General as a takeoff point for a SIL." (*78/05/11*)

Dennis R. Elliss, Cray Research, Boulder, CO 80303: "I have a COPYSF (copy shifted file) implemented on a CRAY-1 written in Pascal using 11 lines of code." ( $* 78 / 08 / 07 *$ )
Larry Ellison, Computer Assisted Bible Study, Willingboro, NJ 08046: "I am serving as coordinator for a group of Bible students who are going to use Pascal on various coordinator for a group of Bible students who are going to u
micro-computers to assist in the study of the Bible." ( $* 78 / 08 / 09 *$ )

John Fitzsimmons, Edina, MN 55436: "It seems that every issue of PUGN has a few pleas for insertions, deletions, or things they don't like about Pascal. Did it ever occur to those of you who complain that the rest of us like the language as it is?" (*78/06/30*)
Lee Fränk, BTI Computer Systems, Cherry Hill, NJ 08002: "... our Pascal is the systems programming language for our new BTI-8000 and all our compilers are written in it." *78/06/16*

Glen Fullmer, Tektronix Inc., Beaverton, OR 97099: "Dear Lord, won't you buy me a new programming language/ My friends all write Pascal/ I must make ammends./ P. S. We could call it 'LACSAP'." (*78/10/31*)

Steven J. Greenfield, Unicorn Systems Company, Los Angeles, CA 90010: "I have been using Pascal for the last six months to write an Assembler designed to generate code for any object computer. Pascal has provided a powerful method of writing a transportable piece of sof tware." (*78/04/25*)

Dale H. Grit, Department of Computer Science, Colorado State University., Ft. Collins, CO 80523: "At CSU, we're using Pascal
course will be Pascal." ( $* 78 / 08 / 10 *$ )

## Here and There With Pascal

Marc Hanson, Hermosa Beach, CA 90254: "... I would appreciate learning about anyone' experiences with running Pascal on either Xerox or Honeywell equipment." ( $* 78 / 05 / 04^{*}$ )

G. Steve Hirst, Iowa City, Iowa 52240: "CONDUIT (a consortium distributing computer-based curriculum materials) is currently investigating including Pascal as a distribution language for new materials." (*78/08/07*)

Claes Hojenberg, University of Uppsala, Sweden: "UDAC is the computer center for the Univ. of Upsala, Sweden's biggest university, and we hope to be able to use \{UCSD\} Pascal for implementing a data-base management system on microcomputers." (*78/10/06*)
K. B. Howard, College of the Sequoias, Visalia, CA 93277: "We're interested in looking into the possibility of using Pascal (in instucting beginning programming course) for
students aiming toward engineering and computer science fields, and are particulary students aiming toward engineering and computer science fields, and are particulary interested in learning of sources for compilers for the language, for PDP-11, HP-3000, and/or Altair 8800 micro if possible." (*78/09/29*)
L. C. Hutchinson, Mentor, 0 H 44060: "... I would appreciate knowing if there are any Modcomp Pascal users..." (*78/05/15*)

Jose I. Icaza, Universidad Autonoma Metroplitana - Azcapotzalœo, Mexico D.F., Mexico: "At this University, we are just starting to use Pascal and giving some optional mini-courses about it. People seem to love the language. Hopefully, soon it will replace FORTRAN as the first language students learn." (*78/10/24*)

Dennis Kalthofer, Philadelphia, PA 19103: "I am starting a workshop in computer science stressing the social aspects of the field. ... I plan to use Pascal as the basis for these systems and any further systems we develop, to organize our programming technique and understanding, and for teaching people about computers and programming in general, as it illustrates many important computer concepts." (*78/07/11*)
Richard H. Karpinski, San Francisco, CA 94114: "Request that software tools' or 'applications' solicit Pascal program modification tools, such as macro generators and programs to make names unique among the first $N$ characters, etc. praise for UCSD system." (*78/04/18*)

Tom Kelly, Downingtown, PA 19335: "With regard to "improvements", "extensions", etc, wish people would engage brain before putting mouth in gear'. My (substantial) work with several Pascal compilers over past year has shown me what a fantastic job Wirth did!" (*78/07/07*)
Neb Lafert, Hewlett-Packard (Schweiz) AG, Geneva, Switzerland: "... we think that a good relationship should be established between our two organizations, enabling us to help every new request for Pascal in our country." ( $* 78 / 09 / 25 *$ )

Jerry LeVan, Dept Math Sc, Eastern Ky Univ, Richmond, KY 40475: "All of our CS majors will be started on Pascal. We are using OMSI's Pascal. I am reasonably happy with the implementation (it will compile and run Pascal-S)." (*78/07/11*)
Stephen A. Locke, Beloit Corporation/Paper Machinery Division, Beloit, WI 53511: "I am interested in Pascal for real-time control of an industrial process... Is there anyone you know working in such a direction?" ( $* 78 / 06 / 05 *$ )

Richard C. Lound, San Francisco, CA 94114: "I am an independent software consultant, primarily in communications systems. My interest in Pascal is in its applicability to
use for generation of specialized message switching and front-end software. $1\left(* 78 / 08 / 02^{*}\right.$ )

Wilf Overgaard, Worldwide Evangelization Crusade, Fort Washington, PA 19034: "Where could I locate a general ledger-bookkeeping program, in Pascal, for non-profit organization? Where can one find a good word processing program in Pascal?" ( $* 78 / 08 / 31 *$ )

Bill Marshall, Sanders Associates Inc., Nashua, NH 03060: "I had hoped to be the first one on my world to implement Pascal on the VAX-11/780, but discovered a group at Univ.

## Here and There With Pascal

Jim McCord，Goleta，CA 93017：：I am acting as the distributor for UCSD Pascal for hobby users of the LSI－11．Cost is $\$ 50$ ，of which $\$ 35$ goes to UCSD for continued work．Other
$\$ 15$ pays for documentation and postage，if user sends me four floppies．（Else I will provide for $\$ 3$ each）．This includes all source code for everything，including the （ $* 78 / 07$（17＊）Anybody interested should get in touch with me（we already have 7 users）．＂ （＊78／07／17＊）

Michael Robert Meissner，University of Minnesota，Minneapolis，MN 55455：＂Everybody talks about portability of programs．This summer I ran into the portability of programmers． Pascal compilers，and have to＇relearn＇Pascal whenever we switch computer systems or compilers．＂（＊78／10／20＊）

Anne Montgomery，Lowry AFB，CO 80230：＂McDonnell Douglas has developed a CMI／CAI system here on Lowry Air Force Base called the Advanced Instuctional System（AIS）．The AIS，as its name implies，is used primarily for technical training．．．．The sys tem currently nanages approximately
$(* 78 / 10 / 16 *)$

Roderick Montgomery，Somerville，NJ 08876：＂I am coordinating distribution of UCSD Pascal to amateurs in the Amateur Computer Group of New Jersey，largest surviving hobbiest club in U．S．September meeting of ACG－NJ will be devoted to Pascal．＂（ $* 78 / 07 / 20 *$ ）
William Moskowitz，The California State University and Colleges，Los Argeles，CA 90036： ＂I might add that Pascal at CSUC has been tremendously successful．During the past twelve months we have had 68,603 accesses and usage continues to grow．＂（ $* 78 / 07 / 17 *$ ）

David Mundie，104－B Oakhurst Circle，Charlottesville VA 22903：＂I would like to correspond with anyone having first－hand experience with the s－100 bus TI 9900 Pascal system being offered by Marinchip Systems．＂（＊78／10／06＊）

John E．Newton，Randolph AFB，TX 78148：＂I am specifically interested in identifying members that have implemented Pascal on Burroughs 6700 hardware．＂（ $* 78 / 07 / 20 *$
Dave Peercy，BDM Corp．，Albuquerque，NM 87106：＂We at BDM are becoming a very interested group of Pascal users．＂（ $* 78 / 08 / 28 *)$
ergi Pokrovsky，USSR Acad．Sci．，Novosibirsk，USSR：＂I hope that S．Pitin of the Moscow Computing Center will shortly report to you on his（not so recent）implementation of Pascal for the BESM－6 computer．＂（＊78／10／31＊）
Darrell Preble，Georgia State University，Atlanta GA 30303：＂GA State Univ．has converted a Pascal compiler from SUNY at Stony Brook．Originally written in XPL，it converted a Pascal compiler from SUNY at Stony Brook．Originally written in
uses either of two monitors to support interactive or batch use．＂$(* 78 / 09 / 05 *)$

David Rosenboom，York University，Downsview，Toronto，Canada M3J 1P3：＂My particular nterest in Pascal is in obtaining or developing a compiler for use on the $16-b i t$ Interdata machine．．．Do you know of anyone who has developed a Pascal system for Interdata 16－bit machines？＂（＊78／09／01＊）
Axel Schreiner，University of Ulm，W－Germany：＂Using（in Ulm）Torstendahl＇s RSX－11 Pascal （love it）and Petersen＇s TR440 Pascal（not quite as stable）in beginner＇s courses．＂ （＊78／06／19＊）
oeseph C．Sharp，Varian，Palo Alto，CA 94303：＂I will introduce Pascal to the North Star Users Group this month．A 30 minute talk is scheduled．＂（＊78／10／30＊）

Robert J．Siegel，Brooklyn，NY 11215：＂Would like to see an article on the relationship of Pascal to ALGOL．＂（＊78／06／23＊）
Seymour Singer，Hughes Aircraft Co．，Fullerton，CA 92634：＂We have installed the SLAC－Stanford Pascal compiler on our twin Amdahl 470 computers．＂（ $* 78 / 07 / 09 *$ ）

Jim Smith，Computer Science Dept．，School of the Ozarks，Pt．Lookout，MO 65726：＂We have recently implemented a Computer Science Department here as the School of the Ozarks，and there is a need to increase the software library in the computer center．We feel that Pascal would be an important language lo present in the curriculum．＂（＊78／09／08＊）
Craig A．Snow，TRW Communications Systems and Services，San Diego，CA 92121：＂We are very interested in using Pascal to implement our future sof tware products．＂（ $* 78 / 05 / 09 *$ ）

James A．Stark，Oakland，CA 94609：＂Interactive Pascal via UNIX is way ahead of a batch compiler on UCSF＇s 370／148 but I have no comparison on routine production jobs on either．＂（＊78／07／17＊）

Ed Thorland／Walt Will，Computer Center，Luther College，Decorah，IA，52101：＂We are still looking for information on an HP3000 implementation of P－code Pascal．Also need documentation of p－code instruction－format and functions．＂（＊78／07／11＊）

P．J．Vanderhoff，Berkel En Rodenrijs，The Netherlands，＂What happened to Stony Brook Pascal release 2？＂（ $* 78 / 10 / 27 *$ ）

Eiiti Wada，Division of Engineering，University of Tokyo Graduate School：＂In my class， all the examples were switched to Pascal since the fall semester of 1972，and the firs Pascal compiler became available in the summer of 1974．Since then at the University of Tokyo，three versions of Pascal compilers have been installed，and all the compilers are
intensively used．＂（ $* 78 / 09 / 08 *)$

Anna Watson，Panama City，FL 32407：＂Very fascinating reading in News－must obtain magnifying glass before I go blind though．＂（＊78／05／15＊）

Anna Watson，Panama City，FL 32407：＂Is there a Pascal for a SEL 32／75？＂（＊78／10／07＊）
John West，Digital Systems Design Group，Atlanta，GA 30327：＂Would like any information 2 （＊78／05／01＊）

James A．Woods，Berkeley，CA 94703：＂What＇s wrong with C？＂（＊78／08／24＊）
PASCAL IN THE NENS
Byte，May，1978：＂Comments on Pascal，Learning How to Program，and Small Systems＂；A short article by Gary A．Ford，Arizona State University，which talks about Pascal s advantages and drawbacks with regards to personal computing．II have used Pascal for at least 95\％ of my own programming and I cannot recommend it too strongly．＂
Byte，August，1978：＂Pascal：A Structurally Strong Language＂；A 6－page article describing Pascal．Polish conversion，and subsequent code generation for a hypothetical micro are listed and explained．

Byte，August，1978：＂In Praise of Pascal＂；A quick survey of Pascal，with descriptions of user－defined scalar types，sets，and pointer type variables．A comparison of a pascal program and a BASIC program to its corresponding Warnier－Orr logic diagram is given．
Byte，August，1978：＂Pascal Versus COBOL＂；Ken Bowles shows how Pascal can be applied to the traditionally COBOL－infested business environment

Byte，August，1978：＂Pascal Versus BASIC＂；A comparison of a program＇MASTERMIND Codebreaker＇written in both BASIC and Pascal．Mastermind is similar to the number guessing game＇BAGELS＇，using colored pegs instead of digits．

Byte，September，1978：＂A＇tiny＇Pascal Compiler，Part 1：The P－code Interpreter＂；The first in a series of articles describing a Pascal compiler written for an 8080．The first talks about parsing，and grammers，etc．Parts of the $P$－code interpreter are listed．

Byte，October，1978：＂A＇tiny＇Pascal Compiler，Part 2：The P－compiler＂；The second part of the previous，this describes the compiler portion．

Byte, November, 1978: The third part of the 'tiny' Pascal series is to be on generating executable 8080 machine code
Computer Week, May 12, 1978: "Pascal- Everybodys Language?"; A short description of what Where, and Why of Pascal. "Pascal is named after the 17th century French philosopher Where, and why of Pascal. ascal. It is not an acronym and is written in lower case."

Computer Weekly, August 24, 1978: "GEC's Pascal"; "A Pascal compiler is being developed by General Electric Computers (GBr) for its 4000 series machines
will be available in 1979.

Computer Weekly, September 9, 1978: "Motorola to offer Pascal on MACS"; "Giving futhe credence to the view that Pascal could become the dominant high level language of language supported on its new microprocessor MACS, due to be unveiled early next year."

Computerworld, April 17, 1978: $I I$ adds Pascal to Mini's Repertoire"; "A Pascal software package said to be suitable for systems applications because its compiler and several software modules are themselves written in Pascal has been introduced by Texas Instruments Inc. for the firm's DS990 packaged disk-based minicomputer systems." 1 year sof tware subscription costs $\$ 1,500$ to $\$ 2,000$.

Computerworld, April 24, 1978: "Growth in Use of Pascal called Revolutionary"; A short report, by Richard Cichellii, mentioning that Pascal is available "for the Zillog, Inc.
Z80 micro to the Cray Research, Inc. Cray-1 supercomputer and for nearly everything in between." Also, it gives the addresses of the PUG and DECUS/Pascal groups.

Computerworld, May. 8, 1978: "Pascal Attractive Anyway"; A Letter to the Editor from Sau. Rosen, Purdue University, "Pascal is a very attractive language. Here and at many other colleges and universities, it is used extensively in computer science and compute engineering courses.

Computerworld, May 15, 1978: "Standard Pascal Compiler Runs on PDP-11's"; A description of Oregon Minicomputer Software, Inc. Pascal compiler, known as OMSI Pascal-1, which generates assembly code that can be assembled and linked with DEC system utilities. RT-11 can support this compiler.
Computerworld, May 22, 1978: "Pascal ready for DG users"; An announcement of Rhintek Inc.'s Pascal compiler for Data General Corp. minicomputers running RDOS. Cost is approximately $\$ 1,000$

Computerworld, May 22, 1978: "Northwest Melds 8085A, Pascal"; "Northwest Microcomputer Systems, Inc. has announced a 'programmers workbench' that reportedly combines the throughput of the 3 MHz Intel Corp. 8085 A and the power of Pascal." "The $85 / \mathrm{P}$ provides sequential Pascal environment, according to the spokesman, including ranalil documentation plus a 90 -day warranty." Cost is about $\$ 7500$.
Computerworld, September 4, 1978: "University Working To Adapt Pascal For MDC-100 Use"; "Programmers here at the University of California are presently under contract to adapt
Pascal for use on the American Microsystems, Inc (AMI) MDC-100 microprocessor development pascal for use on the American Microsystems, Inc (AMI) MDC-100 microprocessor developmen

Computerworld, September 29, 1978: "The Waves of Change", "Implementation languages and the case for Pascal"; one section of the multi-part excerpt of Charles P. Lecht's book, $\frac{\text { The }}{\text { and }}$ Where $\frac{\text { of }}{\text { it }} \frac{\text { Change }}{\text { is being devoted to a background of why Pascal is a successful language, }}$ development languages such as Algol 68, because it is apparent that it was designed for development languages such as Algol 68, because it is
software engineering purposes. (italics in original).

Computerworld, September 25, 1978: "Isam Logic, Disk Space Control Included in Micro-Based Pascal/ $Q^{\prime \prime}$; An announcement of Pascal/ $Q$, which is an enhanced version of Pascal which "includes support for Qsam, Queue's enhanced Isam fille access method, and for automatic disk file storage allocation". Available for $\$ 300$ plus $\$ 19 /$ month for updates from Queue Computer Corp.

Computerworld, October 2, 1978: "DOD Expects Standard Compiler by 1981"; The U. S. Department of Defense's new compiler is planned to be based upon Pascal. There is a plethora of articles on this language (see July Sigplan Notices).

Computerworld, November 20, 1978: "Work on Pascal Progressing"; "A technical committee from the American National Standards Institute (ANSI)] ON Pascal has been approved to work under the X3 committee on computers and information processing. Identified as X3J9, the new groups' initial task is to prepare a proposal for standardization of Pascal and to obtain approval of the proposal..." Justin Walker of the NBS will convene the first
meeting at the CBEMA offices on Tuesday, December 19. "Interested people and organizational representatives are invited to contact Cathy A. Kachurik at Cbema/Standards, izational representatives are invited to
1828 L St. N. W., Washington, DC 20036 .

Computerworld, November 20, 1978: "DOD language named"; "ADA" has been chosen as the name for the forthcoming Department of Defense (DOD) computer programming language. The language was named after the first programmer in history, according to Lt. Col. William A. Whitaker of the DARPA. Ada Augusta, Countess of Lovelace, was one of the few
contemporaries of computing pioneer Charles Babbage who understood his work on calculating machines. ... the first funded compiler, produced for the Army is expected in May 1981."

Computing Europe, September 1978: "Steelman ready next April ..."; More on the DoD's new language. Some background on what has been happening, plus some comments by Edsger Dijkstra, who is a critic of the DoD's plans.

Electronic Design 19, September 13, 1978: "Pascal isn't just one more computer language. It promises to be simple, flexible and fast."; This introduction to the Pascal programming language is the first part of a series, based on ESI's Pascal Instruction Manual. Future parts will deal in detail with Pascal statements, structured data, I/O ood primer to the language. About 5 pages in length.

Electronic Products, July 1978: "As IC it"; bylined by Jerry Metzger. He mentions that several IC houses and minicomputer companies have announced intentions towards using Pascal." But standards need to be established. The time is right to do this with pascal."

Electronic Engineering Times, October 16, 1978: "Pascal Implemented in Code of WD's First Computer Offering"; Pascal has been implemented in the microcode of a new computer from according to the company." "This new system includes a complete Pascal operating systemPascal compiler, BASIC compiler, file manager, screen-oriented editor, debug program and graphics package- all written in \{UCSD\} Pascal." Price is about $\$ 2,500$.

Electronics, October 12, 1978: "Pascal becomes software superstar"; "From the mountain fastness of Switzerland there came 10 years ago a programming language called Pascal. For the first few years of its life it created little stir, but then it began to gain popularity in academia and eventually industry. Today, Pascal is finding its way into rief history, and the current usages of Pascal, from micro's to maxi's and small applications programs to operating systems.

Scientific-Technical Book \& Copy Center, Letter to Andy Mickel; "Pascal is our best seller號
Silicon Gulch Gazette, Volume 3, Number 3: "UCSD Pascal on An S-100 System"; "Dr. Jim pascal on his small computer and the difficulties fnvolved in the project." report on scheduled lectures during The Third West Coast Computer Faire, which took place November 3 rd and 4 th.

From the preceeding: "A Portable Compiler for a Pascal-like Language"; "... will be described by Mark Green. He will treat the problem of program portability. Three solutions to the problem will be presented. As well, a particular piece of portable sof tware developed for the Micro Pascal Compiler will be examined. "

Communications of the ACM, October 1978, back cover: An advertisement for jobs with the Software Technology Company; "develop a compiler for a sophisticated, Pascal-based communications language with real-time multiprocessing features, extensive exceptionhandiling facilites, global data modules, and other state-of-the-art characteristics." "\{Softechs\} compiler was produced on the UNIX system and later moved to RSX-11."

TimeShare, open letter to PUG members: "TSC has adopted Pascal as the primary implementation language for its LSIll based products. ... It is, however, difficult to fin programmers experienced with Pascal
applicants with these qualifications
(plus $2-4$ years exp erience).

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612) 374-1293 or PUG member Eric Hand, 2633 Dupont Ave S., Minneapolis, MN 55408 (612) 377-7387, informed us that if you are tired of cumbersome COBOL, obsolete FORTRAN, or Brontosaurian PL/1, you can show your support for the best general-purpose language now available by
cquir

PASCALINTEACHING
This new section will report on experiences with Pascal used for teaching in computer science. The first report is a nice survey done in Australian Universities by Jan Hext from the University of Sydney. Following that is a report from Japan, and one in a CAI system developed at ETH Zurich. Judy Bishop at the $\begin{aligned} & \text { in } \\ & \text { Ion }\end{aligned}$ contest held for undergraduates. Substantial prizes were given.

## The University of Tasmania

Postal Address: Box 252C, G.P.O., Hobart, Tasmania, Australia 700 Telephone: 23 0561. Cables 'Tasuni' Telex: 58150 UNTAS

Dear Andy,
11th October, 1978

I enclose some information which should be of interest to Pascallers. A friend of mine, Jan Hext from the University of Sydney, has been polling Australian Universities to measure the extent of Pascal's penetration into the teaching area. The sampling is very selective (ie. by membership of PUG!) or would contribute insignificantly. There are exceptions, of course, notably Monash University - I am reliably informed they are switching over in 1979.

## TEACHING PASCAL IN 1979

In order to survey the market for Pascal textbooks in 1979, a questionnaire was sent to the universities and colleges listed in the Pascal Users Group mailing list. Three questions were asked:

1. How many students would be learning Pascal in 1979 ?
2. Would they have learned any other language previously? If so,
which one?
3. What textbook would be recommended?

The answers are summarized in the table below. Allowing for a few selftaught students, etc., the main conclusion is that at least 2500 people in Australia will be learning Pascal in 1979, of whom 1900 will be learning it as their first language.

Also listed below are thirteen textbooks on Pascal which are either available or else in press.

The enthusiasm for Pascal may be reflected in the fact that all of the questionnaires were returned without any extra prompting. I would 1

Replies to questionnaires

| University or <br> Institute | Introductory <br> Students | As a Second Language: <br> students, first language |
| :--- | :--- | :--- |
| Adelaide | 350 | 40, Fortran |
| A.N.U. | 250 | - |
| Melbourne | $200(?)$ | 100, Fortran |
| Newcastle | - | 35, Fortran |
| N.S.W. | 320 | - |
| Queensland | 400 | 100 , Fortran |
| R.M.I.T. | 150 | - |
| S.A.I.T | - | 100, Cobol |
| Sydney | - | 200, Fortran |
| Tasmania | 120 | - |
| W.A. | - | - |
| Wollongong | 1890 | 60, Basic |
| Total |  |  |

## Textbooks

The following textbooks are either introductions to Pascal or more advanced books that make use of Pascal. Reviews of them are cited from the Pascal Newsletter (PN) and the ACM Computing Reviews (CR)

Addyman and Wilson: "A Practical Introduction to Pascal", MacMillan, 1978, 140 pp.
Alagic and Arbib" "The Design of Well-Structured and Correct Programs", Springer, 1978, 292 pp. (PN\#11).
Bowles: "Microcomputer Problem Solving Using Pascal", Springer, 1977, 563 pp . (PN\#11).
Conway, Gries and Zimmerman: "A Primer on Pasca1", Winthrop, 1976, 448 pp. (PN\#12).
Findlay and Watt: "An Introduction to Progranming in Pascal", Pitman, 1978. Grogono: "Programming in Pascal", Addison-WesTey, 1978, 350 pp. (PN\#12). Jensen and Wirth: "Pascal Users Manual and Report", Springer, 1978, 167 pp. Kieburtz: "Structured Programming and Problem Solving with Pascal", Prentice-Ha11, 1977, 320 pp . (PN\#10)

Roh1 and Barrett: "A First Course in Programming in Pascal", Cambridge University Press, in press.

Schneider, Weingart and Perlman: "An Introduction to Programming and Problem Solving with Pascal", Wiley, 1978, (PN\#12).
Webster: "Introduction to Pascal", Heyden, 1976, 129 pp. (PN\#8).

Wirth: "Systematic Programming: An Introduction", Prentice-Hall, 1973.
Wirth: "Algorithms and Data Structures = Programs", Prentice-Hal1, 1976.
NIHON UNIVERSITY
COLLEGE OF INDUSTRIAL TECHNOLOGY

$$
\begin{aligned}
& \text { Izumicho Narashino Shi } \\
& \text { Chiba 275 Janan }
\end{aligned}
$$

A Report from College of Industrial Technology
Nihon University, Japan
Prof.H.Shima feel strongly the fruitfull effect of utilizing the Pascal language in computer science education, and so he utilize that language in his class.
The year Prof.h. of the department of mathematical engineering was attended'to it's seminar. The first semester of 1977, he utilized Pascal for 110 students of junior enroll to the department in computer science class, and all these times they used "Systematic Programming: An Introduction" (Prentice Hall' 71 :-Translated to Japanese Edition) as a text.

Now, in 1978' academic year, on both former term and later term he use mainly Pascal in his class for computer, science education, referring "Algorithm + Data StructPascal in his class for computer science education, referring "Algorithm + Data Stru
ures $=$ Programs" (Prentice Hall'76) and using a text note which Prof.H.Shima himself edited for his junior level students and they belong to the department of mathematical engineering.

Students are served to use concurrent Pascal compiler for their practice and it is implemented by Assistant Prof.J.Ohshima on his laboratory minicomputer (Facom U-mate).

## XS-0

In the Apr/May/Jun 1978 issue of the AEDS Monitor, an article appeared entitled XS-0 "XS-0: A Self-Explanatory School Computer" by J. Nievergelt. The paper was presented at the NAUCAL 1977 Fall Computer Conference in Dearborn, MI. Nievergelt is with the Institut fuer Informatik, ETH Zurich and also with the Department of Computer Science at the University of Illinois. Other people involved in the project are H. P. Frei,
H. Burkhart, Chris Jacobi, B. Pattner, H. Sugaya, B. Weibel, and J. Weydert also of H. Burkhart, Chris Jacobi, B. Pattner, H. Sugaya, B. Weibel, and J. Weydert also of ETH. The project, begun in Fall, 1975, was intended to develop an interactive system that should serve as a self-explanatory school computer so that a user should be able both as an author language and as the programming language for teaching purposes. The hardware consisted of a PDP 11/03 with 28 K words and dual floppys, 2 graphics terminals with TV monitors and 8080 micros with 8 K bytes of RAM. The system software was written in MODULA. The 8080 was programmed in assembler.

Latest News About DOD-1 (ADA or DODD)

- Andy Mickel

As we've told you in previous issues of Pascal News, the U. S. Department of Defense (DOD) has endeavored to procure a common programming language based on Pascal for all "embedded" computer applications--computer systems attached to weaponry. Reliable software should kill people reliably! A series of proposals were drawn up under the alternatively titled "Department of Defense Requirements for High Order Computer Programming Languages." The DOD awarded 4 contracts to 4 software houses from those who had responded to the Ironman specifications in July, 1977. They formulated actua language designs in- documents which are known by colors: BLUE-SofTech; GREEN-Honeywell Bull; RED-Intermetrics; and YELLOW-SRI International.

Basically, the designs consist of Pascal extended for concurrent prgcessed and timedependent ("real-time") programming. Because a projected $\$ 3.0 \times 10$ will be spent each alone has stimulated much manufacturer interest in Pascal over the last two years. We were always worried that this new language (formerly referred to as "DOD-1" and which has now been dubbed "ADA" "-see Pascal in the News--or DODO) would swamp Pascal if it were too similar in form. Manufacturers then simply would not support Pascal but instead supply the new, extended language.
In February, 1978 the DOD narrowed the field to 2 by selecting GREEN and RED for actual implementation efforts. More than 50 groups of academic, military, and industrial onsulted for YELLOW (the least ambitious of the proposals) and Henry Ledgard for GREEN t is reassuring that none other than Edsgar Dijkstra wrote caustic comments which appeared in SIGPLAN Notices: EWD663 in July and EWD659-662 in October. ADA is safely going off the rails, and the threat to the integrity of Pascal is over, l believe. To quote Dijkstra:

BLUE - "unacceptably complex"; GREEN - "the mixture between sense and nonsense
remains baffling"; RED - "both advanced and backward in such an incongruous
manner that 1 am baffled"; YELLOW - "an unsalvagable mess. ${ }^{\text {" }}$
He stated in EWD663:
expected to also quite clear why the new programming language cannot be
expected to be an improvement over Pascal, on which the new language should at all-will be much, much worse than Pascal if they proceed in this fashion.) You cannot improve a design like Pascal significantly by only shifting the
'centre of gravity' of the compromises embodied in it: such shifts never
result in a significant improvement, in the particular case of Pascal it will
be extra hard to achieve any improvement at all, as most of its compromises
have been chosen very wisely..."
Please see Ed Reid's letter in the Open Forum section

## BOOKS AND ARTICLES

Please submit all notices of Pascal books, articles, abstracts, etc. to Rich Stevens at the address below.


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After going through the previous Newsletters I decided to break the Books and Articles section into:

- Articles
- Books
did not include any abstracts with each article reference and only included a comment when I felt one was needed for clarification as to the papers relevance to pug. This should cut down on the size of the sectionalittle. 1 expanded the book section and gave as much information on the book as possible (table of contents when available) as this is the kind of stuff that I look at when initially inspecting a book.
I just received your UCC Computer User's Manual today and am initially very impressed (especially with the introduction to computing).. I'll send more detailed comments shortly.


ARTICLES
Amman, Urs, "Error Recovery in Recursive Descent Parsers", ETH Zurich, Berichte des Instituts fur Informatik, No. 25, May 1978.
Berry, R. E., "Experience with the Pascal P-Compiler", Software - Practice and Experience, Vol. 8, 617-627 (1978).

Burger, Wilhelm F., "Parser Generation for Micro-Computers", Dept. of Computer Sciences, U. of Texas at Austin, TR-77, March 1978.
(* A parser for the language Pascal can be accomodated in less than 4 K of 8-bit bytes *)
Erkio, Hannu and Sajanienu, Jorma and Salava, Autti, "An Implementation of Pascal on the Burroughs B6700', Dept. of Computer Science, U. of Helsinki, Finland Report A-1977-1.

Krouse, Tim, Electronic Design, Vols. 19 thru 23, 1978. (* A continuing series of tutorials on Pascal *)

Lawrence, A. R. and Schofield, D., "SFS - A File System Supporting Pascal Files Design and Implementation', National Physics Laboratory, NPL Report NAC 88 Feb. 1978.
LeBlanc, Richard J., "Extensions to Pascal for Separate Compilation", SIGPLAN Notices, Vol. 13, No. 9, Sept. 1978.

Lecarme, Olivier and Peyrolle-Thomas, Marie-Claude, "Self-compiling Compilers: An Appraisal of their Implementation and Portability", Software - Practice and Experience, Vol. 8, 149-170 (1978).
(* The study is centered around a specific case, the programming language Pascal and its many compilers *)

Marlin, Chris D., "A Model for Data Control in the Programming Language Pascal" Proceedings of the Australian Colleges of Advanced Education Computing Conference, Aug. 1977, A. K. Duncan (Ed.), pp. 293-306. Available from author at Dept. of Computing Science, U. of Adelaide, Adelaide, South Australia 5001. Marlin, Chris D.,
Pascal, 'ISoftware - Practice and Experience, to appear. Also available from the author, see above.

Here is the Books and Articles section for \#13. Thank the world for self correcting typewriters. promise to have thingsbetter organized so that my secretary can do the typing for $\# 14$.

Mohilner, Patricia J., "Prettyprinting Pascal Programs", SIGPLAN Notices, Vol. 13, No. 7, July 1978.

Neal, David and Wallentine, Virgil, "Experiences with the Portability of Concurrent Pascal", Software - Practice and Experience, Vol. 8, 341-353 (1978).
Posa, John G.,"Pascal Becomes Software Superstar", Electronics, Oct. 12, 1978.
Posa, John G., "Microcomputer Made for Pascal", Electronics, Oct. 12, 1978.
Pratt, Terrence W., "Control Computations and the Design of Loop Control Mar. 1978.
(* Examples drawn from a Pascal Compiler *)
Sale. A. H. J., "Strings and the Sequence Abstraction in Pascal", Dept. of Information Science, U. of Tasmania.

Sale, A. H. J., "Stylistics in Languages with Compound Statements", Australian Computer Journal, Vol. 10, No. 2, May 1978.

Shrivastava, S. K., "Sequential Pascal with Recovery Blocks", Software Practice and Experience, Vol. 8, 177-185 (1978).
Tennent, R. D., "Another Look at Type Compatability in Pascal", Software Practice and Experience, Vol. 8, 429-437 (1978).
BOOKS
$\frac{\text { PASCAL: An Introduction to Methodical Programming by Bill Findlay and David Watt }}{\text { (U. of Glascow, Computing Science Dept.). Computer Science Press, } 306 \mathrm{pp} \text {; }}$ (U. of Glascow, Computing Science Dept.). Computer Scie
UK Edition by Pitman International Text, 1978 ( $£ 4.95$ ).

The book does not assume previous knowledge of computing, nor of advanced mathematics. Emphasis is placed on programming principles, good style and a methodical approach to program development. The technique of stepwis major chapters are exclusively devoted to programming methodology. The first is placed early enough to encourage good practice from the start. It includes sections on choosing refinements, testing and correcting and documentation. The second, at the end of the book, draws all the materia together in two realistic case studies. Since the whole language is covered, the book may be of value to those who wish to learn something of the modern concepts of program structure and data structure, even if they Part 1: First Steps in Programming

Computers and programming; data and data types; the INTEGER type;
HILE and IF; methodical programming, Case Study 1.
Part 2: More Data Types
CHAR, enumerated types, subranges; REAL; arrays
More Control Structures
Par't 4: Subprograms
Functions; procedures, parameter passing, procedures and program structure; advanced uses of procedures.
Part 5: More data structures
Records; strings; files; sets; pointers.
Part 6: Programming Methodology
Case studies 2 and 3 , general principles
Appendices
Collected syntax diagrams; reserved words and special symbols; predeclared entities; legible input and output; character sets. polected exercises.
(* Author's information *)

Programming via Pascal by J. S. Rohl and Barrett (U. of Western Australia), Cambridge University Press, in press.
(* Anybody have any more information on this text ? *) $\frac{\text { A Practical Introduction to Pascal by I. R. Wilson and A. M. Addyman, Springer- }}{\text { Verlag New York, 1978, } 145 \mathrm{pp} \text {. ( } \$ 7.90 \text { ): MacMillan, London, 1978, ( } \mathrm{F} 3.50 \text { ). }}$

Suitable for beginners and experienced programmers who wish to learn the complete Pascal language, this concise introduction includes

- Syntax diagrams and complete examples illustrating each feature of the language;
- Simple problems introducing control constructs, expressions and the use of procedures;
A discussion of the concept of data type, followed by a complete
description of the data structure facilities of Pascal;
- An analysis of more advanced procedures and dynamic data structures;
- Over sixty programs

Contents:
Introduction. The form of a program and basic calculations. Basic control constructs. Variables, constants and expressions. An introduction to input and output. An introduction to procedures and functions. Data types. 1, 2, 3 and 4: Set, array sequential file. Elementary structured types and functions. Dynamic data structures.
(* From publishers information *)
(: See below for review *)
The Design of Well-Structured and Correct Programs by S. Alagic and M. A. Arbib, Springer-Verlag New York, 1978, 292pp. (\$12.80)
(We are
BOOK REVIEW
Programming in PASCAL by Peter Grogono
Addison-Wesley, Reading, Mass., 1978, $357 \mathrm{pp} ., \$ 9.95$.

Finally, an easy to read, lucid description of Pascal.
This book is described in its preface as being suitable for an introductory programming course and in addition it should who wants to learn Pascal.

The author made a point to cover the entire language and this is one of the book's strongest points. (One of the other texts on Pascal, A Primer on Pascal by Conway, Gries and functions, pointers, records entire language, omitting sets, a good description of a specific implementation (the zürich CDC system) and this will help one appreciate the implementation of the abstract language on a specific computer.

Another strong point of the book is that it is not just a text on writing programs in Pascal, rather it is a text on programming. The inclusion of a chapter on program design and an appendix on program standards are a welcome addition to any language description, especially if the book is to be used for an introductory text. The data structures covered include linked lists and trees.

The examples used in the text are excellent and well thought out. Wirth's technique of stepwise refinement is provided to show that "Pascal, with a relively small number
of basic constructs, can nevertheless be used effectively to solve problems outside the domain of academic programming."

There are very few complaints that I have with this book Each chapter is followed by a group of exercises (solutions are not provided) and some indication as to the relative difficulty of each exercise would be helpful. There are relatively few typographical errors.

All in all the book is excellent and a long awaited addition to the Pascal literature.
W. Richard Stevens

## BOOK REVIEW

A Practical Introduction to Pascal by 1. R. Wilson and A. M. Addyman
Springer-Verlag, New York, 1978,145 pp., $\$ 7.90$. ISBN 0-387-91136-7
This book admirably fulfills the promise of its title - it gives a concise, well-organized tutorial on how to write programs in Pascal. The complete language is presented in fourteen short chapters. Particularly notable is the attention paid to the data structuring facilities of Pascal: fully six of the
chapters deal directly with data structures.

After an introduction in Chapter 1, the basic structure of a complete Pascal program is shown in Chapter 2. Chapter 3 describes the control structures available in the language and gives advice on their use (including obligatory warnings about GOTO's). Chapters 4 and 5 discuss variables, constants, expressions, and input/output. Chapter 6, "An Introduction to Procedures and
Functions", is especially good: the appearance at this point in the course of Functions', is especially good: the appearance at this point in the course of the presentation of these concepts is well-motivated and natural. Also, Pascal's including files and record variants. Procedures and functions are revisited in Chapter 13 to show recursion and in Chapter 14 pointers are introduced in the context of "dynamic data structures". Each chapter is followed by suitable sets of exercises(easy) and problems (hard). There are four appendices: the completesyntax, delimiter words, answers to exercises and suggestions for solutions to the problems, and a note about the Pascal User's Group.

The pace of the presentation is even and well-motivated. New syntactic forms are introduced with simplified syntax charts andexamples and their semantics the new Pascal user is the printing of programs as they might actually be ifsted alongwith those ugly digraphs " $(* 11$ and $1 \%$ )". Keywords are, however printed in boldface. Example programs are developed in good style - stepwise refinement and top-down design are advocated and used.

The book is not above some minor criticism: some references are too broad ".... readers are referred to Coleman (1978), Dahl et al. (1972) and Aho et al. (1974)." appears on page 69; refinement of program steps proceeds from comments
expressed in Pascal comments later in the book (page 60 ff) but by lowercase fragments earlier (page 19, ff); there are a few misprints. Also, some of the exercises and example programs would be easier to understand if samples of their input and output were presented.

In summary, the book is a welcome addition to the Pascal literature. It is physically attractive and provides an excellent introduction to the language for beginning and experienced programmers alike.
> R. Warren Johnson

> Department of Mathematics and
> St. Cloud State University

St. Cloud, Minnesota

## CONFERENCES

We received recently, the latest Bulletin de Liaison du Sous-Groupe Pascal no. 4
from Olivier Lecarme in France. He of course heads the French AFCET Pascal Group. This issue of the Bulletin was 125 pages long and is beginning to look like an issue of Pascal News! It contained an editorial, bibliography, list of Pascal implementations, Working Group on Pascal Extensions supplied by Olivier, and if we only had the time it would be the quickest thing to do to translate and print in the Open Forum section. The contributions in the articles section are:

- Pointers: False Problems and Real Insufficiencies by M. Gauthier.
- A Graphic Extension for Pascal by N. \& D. Thalmann.
- An Aspect of TSIMONE: A Version for Pascal Program Profiles by D. Renault.
- Where is the Standardization of Pascal? by O. Lecarme.
- A Comparison and Contrast between Concurrent Pascal and Modula by R. Rousseau.
- An Efficient Method of Controlling Type Unions by Nguyen Van Lu.

An ACM/SIGPLAN - Pascal User's Group sessions is being held at ACM ' 78 in Washington DC. See below.

The Australian Computer Science Conference will hold a workshop on Pascal. The conference is scheduled for February 1 and 2 in Hobart, Tasmania. Arthur Sale, of course is the host and is currently serving as vice-president of the Australian Computer Society. This is the second year for this conference. It was successfully launched under the name Australian Universities Computer Science Conference which was enthusiastically receive last year.
Finally the University of California at San Diego (UCSD) Summer Workshop on Extensions was held this last July and has been reported on by Richard Cichelli below. I was The major results of the Workshop were to get together a variety of computer manufacturers with some dyed-in-the-wool Pascalers. The Workshop rebuffed nearly all proposed extensions except those referred to the International Working Group on Pascal Extensions (such as otherwise for a case statement--see Open Forum under Standards). The members of the Workshop including the more than 15 manufacturers unanimously endorsed a motion to support the speedy adoption of the BSI/ISO Pascal Standard under development by Tony Addyman and his team...see Open Forum

SPECIAL INTEREST GROUP ON PROGRAMMING LANGUAGES

Association for Computing Machinery
1133 Avenue of the Americas
New York, NY 1003
(212) $265-6300$

REPLYTO: Mail Stop 125A
Research Center Hampton, VA 23665

July 24, 1978

Dear Andy,
n informal evening session devoted to PASCAL will be held at the 1978 ACM conference which will take place December $4-6$, 1978, in Washington, D.C this session is to allow all conference attendees who are interested in PASCAL to get together and interact.

This is not a technical session in the usual sense. However, in order to convey the most information, it will consist, at least in part, of a series of short presentations (i.e., approximately 10 minutes) on PASCAL related language and its software; e.g., experience with PASCAL, tools for PASCAL programing, implementations, etc. Anybody who is planning to attend ACM 178 and who is interested in making a presentation should send a short description of what they will discuss by October 1 to:

> John C. Knight Mail Stop 125A NASA Langley Research Center Hampton, Virginia 23665

Presenters will be informed of their selection by November 1 .
The purpose of requesting descriptions is not to perform any refereeing or technical judgment, but merely to allow a balanced program to be prepared for the limited time available.

## Sincerely,


-John C. Knight
tative
978 ACM Conference Program Comittee

## THE UCSD PASCAL WORKSHOP

by

> Richard J. Cichelli ANPA/RI
> Lehigh University

This is a personal report of my experiences with the UCSD Pascal Wowkshop held by Dr. Kenneth Bowles at the University of California workshop, and in no way should this report be considered a report from the workshop participants as a whole.

In May of this year, I received a letter from Ken Bowles inviting me to attend his planned workshop, the purpose of which was to "standardize Pascal on small machines previously at the ACM 1977 Conference. After seeing the UCSD Pascal system in action, I was convinced that it was excellent technology and held great promise for both educational and commercial applications. I reviewed Ken's book, Microcomputer ProblemSolving in Pascal, for PN \#ll and sent a pre-publications copy of that review to Byte Magazine. Upon receiving this review, Carl Helmers
began his own interaction with Ken concerning the UCSD system.

At ACM ${ }^{1} 77$ Ken talked to me about the language changes that he felt small systems required. He spoke of the lack of viable Standards activities within the Pascal Users Group and his willingness to organize a Standards workshop. I suggested to Ken at that time that Standards were something that the Users Group would soon be more involved with and that his help on a PUG Standards Committee would be welcome. I was truly surprised and chagrined to hear of Ken's organizing his own Pascal

Workshop. I have never felt that the precision of expression and depth of understanding necessary for Standard-related activity was the type of thing done well by implementation-oriented individuals. Even more important was the fact that effective international Standards activities had already been initiated in PUG under the direction of Tony Addyman and an Extensions Working Group has been formed chaired by SteengaardMadsen. Dr. Wirth was helping this group with their activities.

On June I, I sent a letter to Ken expressing my concern about the UCSD project. The text of the letter is as follows:
"I have given careful thought to your invitation to participate in the UCSD workshop. ANPA was a member of 27 associations that participated in the acceptance of the 1966 ANSI-FORIRAN Standard. We consider our endorsement of programming language standards of
great importance to our 1200 newspaper members and are sure that
Pascal will have a major impact on future newspaper computer systems Unfortunately, no matter how well meaning your efforts towards
standardized extensions are, we believe the appropriate review and
standardized extensions are, we belleve the appropriate review and
Group. We would welcome your initiative in being part of a PUG
Standards Committee but neither ANPA nor I will support or endorse any self-proclaimed UCSD Pascal modification adventure."

Copies of the letter were forwarded to the Standards Committee, the Working Group, and Andy Mickel. My primary concern with the UCSD effort was that any extensions agreed to by the UCSD froup would become a Standard Pascal. I very much felt that most of the UCSD deviations from the Standard were simply inappropriate. I was sure that most, if not all, of the UCSD language modifications would be rejected both by the Working Group and the Standards Committee. I firmly believe that the UCSD interactive systems feature good engineering. However, like most new implementations derived from the Zurich produced P4 system, the UCSD Pascal fails to implement important parts of the Standard and has extra goodies implemented in ways inconsistent with either the Standard,
or worse, the recommended extension technique.

Upon returning from a business trip, I found that Ken had placed an urgent call to me. I returned his call and spent 2 hours talking with project. Andy and I shared similar reservations. During the conversation, Ken lnvited me to attend the workshop as a PUG representative instead of as a member of a contributing organization. I said that I would give consideration to this idea.

During the month of June I had many conversations with Andy and Other potential workshop attendees. Upon receiving a document titled consisted of more than 75 items, I was even more concerned. In my oplnion, adoption of changes proposed in this checklist would effectively rape the pascal standard. My primary hope at that point was that no one would want to go to the UCSD workshop. Bob Dietrich of Tektronix made a number of telephone calls to me indicating first, that Tektronix was interested in participating in such a workshop and second, that he felt other potential workshop participants felt as I did about the Standard and about these UCSD extensions. He felt that if the issues were properly dealt with, it was likely that the consensus of the workshop would be to reject almost all of the proposed extensions. In later conversations with Ken, he himself also assured me that rejection of ill-conceived extensions would be an important activity of the workshop By this time it was clear that there would be a number of participants in the workshop and that it was important that those workshop participants
who were responsible for corporate implementations of Fascal hear
arguments in favor of adhering to the Standard. Andy. informed me that because of prior committments he could not attend the workshop to represent the Users Group. He asked me to do so in his stead. Also charged with a similar mission were Jim Miner, Arthur Sale, and Bob
Johnson. Since Andy, Bob, and I founded the Pascal Users Group, we ohnson. Since Andy, Bob, and I founded the Pascal Users Group, we hoped

At this point I accepted Ken's invitation to attend the workshop. I also agreed to attend a pre-workshop meeting of like-minded individuals that was the brainchild of Bob Dietrich. Of particular help in formulat by Mike Ball.

The week at the conference was one of the most interesting and challenging of my computer science career. By the sunday meeting we were all found out where the UCBD group stood as far as the extensions these ill-conceived ideas by actually implementing them within the UCSD software system. Shortly after Ken's initial address to the more than 50 participants of the workshop, a number of participants suggested that the overall goals for the workshop be clearly laid out before specific consideration of the checklist items began. A number of views in addition to Ken's were presented on this topic and I was asked to speak on this as a representative of the Pascal Users Group ". Most of
the workshop participants were chosen by Ken because they were members of Pascal implementation teams at various large companies. These peopl are used to identifying problems and developing solutions. I am sure they did not welcome hearing from me that I believed they should act only in an advisory capacity and defer final evaluation to a standards Conmittee within PUG. I am sorry that I don't have a complete transcript of the extemporaneous talk I gave addressing this issue, but the most important point that I tried to make was that ad-hoc solutions to perceived problems with Pascal were to be preferred to hastily the group that if they assured of extensions, the changed language would be a competitor to Standard Pascal, much to the detriment of the user community.

One of the problems in giving this talk was that Ken asked me to present the issues not in the general framework that $I$ just outined but instead as an item by item review of "how do you do ' $x$ '?" (where ' $x$ ' might be direct access files, overlays, complex numbers, strings, etc.) My general statements included a suggestion that the only types of consistent with the design goals of Pascal, and 2) add a facility not implementable in Standard Pascal. For example, in talking about segments and overlays I suggested that such concepts had nothing to do with the problem solved by an algorithm but only with how a compiler translated the algorithm expressed as a program into executable code for a particular operating system. I suggested that if it was necessary for the compiler to know about overlays, then this information should be incorporated in compiler directive comments. (Pascal-6000 needs no such compiler of complex numbers, that they are easily created within the standard mechanisms of the language. I also noted that direct access files ar being considered by the European Working Group. I also mentioned that at Lehigh University we have used direct access files extensively and do so by calling external library routines. Since, at Lehigh, more than four different systems of direct access file support are utilized by Pascal programmers, I suggested that reasonable men would differ as to what constituted a good set of primitive functions for accessing and no consensus about them exists, no standard should be imposed. By not creating a standard for such an item, experimentation is encouraged From this experimentation better solutions can be derived.

During the next three days we broke into subcommittees to consider checklist items one by one. It was Ken's idea that subcommittee sessions would be recorded and "where consensus was reached on an item a consensus position would be prepared". Each subcommittee had one or more UCSD the group's deliberation. A few of these individuals acted as monjtors on their subcommittees.

I worked with what was called the Expressions Group. Our approach was more formalized than some of the other groups. In addition to myself the members of the subdommittee were Terry Miller (moderator), Steven Dum of Tektronix (recording secretary), Ruth Richart of Burroughs, Skip Stritter of MOTOROLA, and Don Baccus of OMSI. We began by considering each of our 15 topics one by one. For each topic we first stated exactly what our recommendations were and then we presented our easoning that went into the recommation item 3.2 on our recommendation is
3.2 We recommend that AND and OR should be left as defined. I.e. the implementor may choose short circuit or complete evaluation user beware!
Short circuit AND and OR (CAND and COR) can be programmed around in existing Pascal. They are a minor extension. The (size, introducing features, etc.) does not justify the benefit

We firmly reject the concept of introducing complete evaluation operators such as LAND or LOR.

On item 3.4 -- "provide for exponentiation" -- we made the following recommendations:
3.4 We recommend rejection of exponentiation as an infix operator or standard function.

It is possible to provide a predefined function POWER or to write it as a Pascal function with the parameters defined as

## FUNCTION POWER (A,B:REAL):REAL

We felt that it was not necessary to add a function to ralse an integer to an integer power as most usage of exponentiation an integer to an integer power as most us

I have the highest regard for the people that Ken recruited to participate in his workshop. Ruth Richart, for example, is a principal implementor of a new systems language that Burroughs is using. This
language is modeled after Pascal. On item 3.2 (the short circuit AND
and OR), she pointed out that on Burroughs machines the short circuit and OR), she pointed out that on Burroughs machines the short circuit evaluation is significantly less efficient than full evaluation in most (and thus super expression are exceedingly efficient on stack operation anditional branches. We concluded that it was important that the mplementor of a Pascal compiler be given the freedom to choose the optimal evaluation technique.

As we worked on each item, we followed Wirth's suggested procedure for considering extensions. First, we introduced the extension in a utorial lashion to the subgroup. Then we showed how the extension would be used in practical programming. Then we discussed its relationship interesting to note that in the AND/OR controversy the UCSD supplied example program was clearly not of the best design. After exploring the
issue for some time, it became clear that short circuit evaluation was most often used in an attempt to sneak past undefined conditions. This led actually to a suboptimal or less clear presentation of the algorithm. It was this kind of discussion that gave us confidence in our recommendations.

At the conclusion of the day's meeting, Steve Dum took our carefully worded notes and typed them into one of the UCSD Terak systems. A little quick editing and we had line printer copy of the day's discussion. The next day we made multiple copies of our preliminary statements on the 15 did not All members of our subgroup were chagrined that the other conclusion of the general meeting on Tuesday, all subgroups were directed to go back and produce concise, well-worded descriptions of their recommendations and deliberations. And the Expression Group was were meeting before the workshop sessions began and after the workshop ended each evening. We were attempting to formulate an appropriate policy statement for PUG which would guarantee that what constituted Pascal was defined by the PUG membership. In this effort we sorely missed having Andy with us to help formulate policy.

By Thursday the work of the subgroups neared completion. Also number of individuals in the workshop indicated interest in working thru PUG on implementations and standards. The following posilion paper to have his name included on the paper:

## PUG Working Position

(1) In October PUG will publish a proposed constitution. Upon acceptance of the constitution by the PUG membership, election of officers will take place. It is hoped that by January 1, 1979, a formal governing structure for PUG will be established.
(2.) A draft of the ISO Pascal Standard will be published by the end of 1978 for member reaction.
(3) An implementation subgroup will be formed to coordinate the enhancement and distribution of portable compilers and to facilitate correspondence among implementors. A new section of Pascal News will inform the membership of these activities.
(4) A standards subgroup will be formed. It will distribute (for a reasonable fee) a Validation Suite. An Incomplete version of the f the Alrol 60 Validation Suite) and distributed by R. Cichelli will be available during September, 1978.
(5) Actual proposals from the International Working Group will appear in Pascal News. The first will be in October, 1978.
PUG aid to the UCSD Workshop
(1) Pascal News will publish (subject to length constraints) a report of the UCSD Workshop and will help to distribute the full Workshop report.
(2) Pascal News will publish a new section on solutions of and commentary about significant programming problems which may be outside the scope of the Pascal Standard.
Andy Mickel
Richard J. Cichelli
Robert Warren Johnson

James F. Miner
Richard J. Cichelli
Arthur H. J. Sale
July 13, 1978

Jim Miner presented the PUG working position paper to the workshop and
It is my opinion that the result of the first week of the UCSD orkshop was to strengthen the Pascal Standard and to reaffirm the pre-emminence of PUG with regards to Pascal. One of the most important factors of that week was the acceptance by all workshop participants of the following "agreement in principle":

At the time the workshop convened, two major activities with
respect to the definition of the language Pascal were already
underway. The International Standards Organization had begun working on a complete definition of the Pascal language in light of the
shortcomings of the Jensen and Wirth document. A Working Group focused around Steensgaard-Madsen had begun working on extensions
to the Pascal Language aimed at correcting a few well-known
deficiencies in the language. In light of these activities the
workshop assumed as its primary goal to address well-defined,
consistent, application-oriented extension sets and agreed to pass to the other two bodies such recommendations and information deemed appropriate to their work.

The workshop recognized the existence of possible modifications to the Pascal Language which, due to the impact throughout the language, would de-facto create a new language and decided not to act on these modifications at this time.
In order to achieve the purposes stated above the workshop has resolved to:
I. Publish and distribute the Proceedings of the workshop. In articular the Proceding will be forwarded to TSO the Pase I Users Group, and the Steensgaard-Madsen Working Group
11. Organize a structure which will permit the orderly continuation of the work begun at the meeting in San Diego.
III. Provide a mechanism to reinforce the importance of Standard Pascal by agreeing that all compilers purporting to support the following statement in the source code and all documentation
"The language --(1)-- supported by this compiler contains th language Pascal, as defined in --(2)--, as a subset with the ollowing exceptions
(a) features not implemented
--(3)-- -- refer to page --
(b) features implemented which deviate from the Standard format
$\qquad$
Notes:
(1) insert the name of the dialect
(2) insert "the Jensen and Wirth User Manual and Report" or "the ISO draft standard" or "ANSI Pascal standard" as appropriate
3) A brief statement plus reference to more detailed information will suffice. The list should be as complete as possible.


## Review of Pascal News 9/10, 11, and 12

Backissues of Pascal News $9 / 10,11$, and 12 are still available, and will be for the forseeable future. Therefore I would like to urge all new members to consider obtaining them so that you will be better oriented to events in our recent past. Issues 1-8 are in detail in issue 11; 1-8 are briefly described in issue $9 / 10$.

If you want to know generally what is important, then issue $9 / 10$ contains the base roster for PUG, and a complete survey of Implementations. It also contains the last wild proposals to extend Pascal, and the terrific article on type compatibility by
Pierre Desjardins. An errata to old printings of Pascal User Manual and Report is in \#11 ssue 12 contains our first applications sections with two important software tools:
COMPARE and a pair of programs for Performance Measurement.
All three issues contain important information about Pascal standards.
$\frac{\text { Pascal News } \# 9 / 10 \text { (combined issue), September, } 1977 \text {, Pascal User's Group, University of }}{\text { Minnesota Computer Center, } 220 \text { pages (114 numbered pages), edited by Andy Mickel }}$
Editor's Contribution: Pascal jobs, a list of computer companies using Pascal, Pascal on personal computers, current information on the status of PUG and Pascal News: Pascal Newsletter to Pascal News, new policies, back issues, growth in membership, and PUG finances.

Here and There: Tidbits ( 9 pages), reports from German and French Pascal conferences Books and Articles classified by applications, languages, textbooks, and implementations; Bibliography of 68 entries; past issues of Pascal Newsletter (1-8); PUG finances for 1976-1977; Roster 77/09/09.
Articles:
"Pascal at Sydney University"
-Tony Gerber and Carroll Morgan
A description of implemented (proven) extensions and changes to the CDC-6000 pascal compiler in use at Uni of Sydney. These include operating system interface dily to read strings, read and write user-defined scalar types, case statement extensions, and two machine-dependent extensions. The conclusion states that thes changes to the compiler have not detracted from the overall efficiency of the compiler, and that 2 -year's use has vindicated the inclusion of these extensions.]
"Disposing of Dispose"
Stephen Wagstaff
[An argument for an automatic garbage collection system for dynamic variables in Pascal is made, thus obviating the need for, and the risks associated with, usercontrolled de-allocation (e.g. DISPOSE). Complete protection from "dangling" pinters may be obtained.
"What is a Textfile?"
ill Price
The definition of the pre-defined type Text in Pascal as File of char is in error and because of this lapse, complex special-case notions are introduced as textfile notion is proposed.]
"Generic Routines and Variable Types in Pascal"

- B. Austermuehl and H.-J. Hoffmann
[Generic routines and variable types, as introduced in ELI are a means to postpone the binding time of routines and data. An examination is given of what degree type checking requirement. The conclusion is made that generic routines fit into Pascal, while variable types have to subject to strong restrictions. Variable types may only be used in connection with a special syntactic form.]

Open Forum:
77/05/10 Arthur Sale to Andy Mickel: [Australasian distribution Centre, CDC-bias: files, program heading, Burroughs 6700 implementation on 7700,6800 , etc.]
77/05/24 Tony Gerber to Andy Micke1: [PUGN distribution to Australia, why haven't you
77/06/01 Richard Cichelli to Andy Mickel: [Each issue of PUGN better, software tools,
77/06/16 Mike incations section in PUGN.] extensions to standard Pascal, proposed standard for editing format and distribution of Pascal software tools and programs.]
77/06/16 Peter Grogono to Andy Mickel: [standardizing Pascal--preserve its simplicity, change to Read procedure for error recovery, especially for interactive programs.]
77/06/24 Wally Wedel to Andy Micke1: [CDC-6000 and DEC-10 Pascal at the Univ. of Texas, standards via X3 and experience from X3J committee.]

正 standard Pascal. Distribution at Colorado is now running smoothly.] services,
using Nage1's DEC-10 Pascal compiler.]
77/07/29 Stephen Soule to Andy Mickel: [Pascal competing with FORTRAN: variable-init ialization, own variables, flexible array parameters, textfiles and variant records in formatting.]
Special Topic: Micro/Personal Computers and Pascal
77/07/08 David Mundie to Andy Mickel: [Zilog rumor about Pascal machine, letters to personal computer journals, game programs in Pascal, like variant records.] 77/06/27 Larry Press to Andy Mickel: [Would like to publish work from PUG members in
77/09/01 Maria Lindsay to Andy Micke1: [Microcomputer library and resource center in
Madison Wisconsin very interested in Pascal materials.]
77/08/24 Jim Merritt to Andy Mickel: [disagree about pressing supposed advantage on micro computers. UCSD Pascal project may hold future hope, UNIX Pascal information.] Pascal. Pascal an excellent choice to succeed BASIC]

Special Topic: Standards
27/08/ Introduction
77/08/09 D. G. Burnett-Hall to Andy Mickel: Another Attention List.
Implementation Notes: Checklist, General Information, Software Tools, Portable Pascals Pascal-P, Pascal Trunk, Pascal J; Pascal Variants: Pascal S, Concurrent Pascal Modula; Feature Implementation Notes: Set of Char, the For statement, Else case, var parameters, Amtahl 470, B1700, B3700/4700, B5700, B6700-7700, CDC Cyber 18 2550, CDC3200, CDC3300, CDC3600, CDC 6000/Cyber 70,170, CDC7600/Cyber 76, CDC Omega 480, CDC Star-100, CII Iris 50, CII 10070, Iris 80, Computer Automation LSI-2, Cray-1, Data General Eclipse/Nova, DEC PDP-8, PDP-11, DEC-10/20, Dietz Mincal 621, Foxboro Fox 1, Fujitsu Facom 230, Harris/4, Heathkit H-11, Hewlett-Packard 21MX, 2100 , 3000, Hitachi Hitac 8000, Honeywell H316, Level 66, IBM Series 1, 360/370, 1130, ICL 1900, ICL 2900, Intel 8080, 8080a, Interdata 7/16, Interdata 7/32,8/32, ITEL AS/4, AS/5, Kardios Duo 70, Mitsubishi Melcom 7700, MITS Altair 680b,

MOS Technology 6502, Motorola 6800, Nanodata QM-1, NCR Century 200,
Norsk Data Nord 10, Prime P-400, Univac 1100, Univac V-70, Xerox Sigma 6, 9, Xerox Sigma 7, Zilog Z-80.
Pascal News \#11, February, 1978, Pascal User's Group, University of Minnesota Computer Center, 202 pages (106 numbered pages), edited by Andy Mickel.

Editor's Contribution: Addenda on list of companies using Pascal. Itemization of costs from PUG(UK) distribution center.

Here and There: Pascal jobs, Help wanted for numerical library project, Tidbits (7 pages), Evolution of PUG dog, Pascal in the News, DOD-1 report, reports from German ACM Pascal meeting and ACM ' 77 Pascal session in Seattle, Books and Articles including Applications, Implementations, Languages, and Textbooks; Concurrent Pascal
literature, documents obtainable from the University of Colorado Pascal distribution
Errata to Pascal User Manual and Report Second edition. Detailed review center. Errata to Pascal User Manual and Report Second edition.

Articles:
Type Compatibility Checking in Pascal Compilers"
Pierre Desjardins
[It is imperative we clearly set down the semantics of type compatibility, for structured variables in the programming language Pascal. The matter is urgent since the lack of an explicit set of rules in that sense has already given rise
to some incompatibilities resulting from the use of different Pascal compilers. to some incompatibilities resulting from the use of different Pascal compilers. On the basis of how a compiler implements type compatibility checking, we can which will react differently to particular cases involving operations on structured variables. It is of course clear that such a conflict must not be allowed to continue, and in that sense I will try to explain how the two classes of compilers came into being and also present the reader with a few examples to display the consequences.]
"A Novel Approach to Compiler Design"
James Q. Arnold
[A sarcastic appraisal of the Honeywell Level 66 compiler implemented by the University of Waterloo. Its poor realization is examined with respect to program program correctness, and user interface.]
"Status of UCSD Pascal Project"
Kenneth L. Bowles description is given of the project which developed the LSI-11 Pascal implementation at UCSD. The project was motivated by teaching interests at the university and has evolved into research and development interests centering on configuration, 8080 and $\mathrm{z}-80$ versions, Pascal extensions and alterations, Introductory Pascal course and textbook, a "Tele-mail" user support facility, and forthcoming improvements.]
"Suggestions for Pascal Implementations"
Willett Kempton
[A user's point of view is presented on features encountered in 3 Pascal implementations. Conditional debugging code, a better cross-reference, flagging nonstandard constructs, implementation of UNPACK, PACK, and LINELIMIT, conversation strings with blanks, and more predefined constants like MAXINT are examined.]
"Suggested Extensions to Pascal"
Robert A. Fraley
[A number of extensions and modifications to Pascal are suggested. It is the author's belief that Pascal as it stands, cannot compete successfully with more complete languages in production environments and over wide ranges of application Some of them are optionally available in the UBC/IBM 370 Pascal compiler.]
"Adapting Pascal for the PDP 11/45"
D. D. Miller
[A description and adaptation is given of the University of Illinois Pascal student compiler for a PDP 11/20, to a production compiler on an 11/45. We will discuss, ) the extensions to the language which were hecessary to communicate between and source update and reformatting, and c) how we introduced Pascal into an existing software system and to MACRO programmers.]
"Pascal: Standards and Extensions"
Chris Bishop
Comments are given on the current standards/extensions argument, and to suggest include: array
read, repeat and case statement changes, inverse to ord, different treatment of file variables, and I/O and textfiles, addition of exponentiation operator.

## Open Forum:

77/11/09 Helmut Weber to Andy Mickel: [CDC-6000 Pascal inquiries]
77/10/28 Barbara Kidman to Andy Mickel: [Pascal teaching at the University of Adelaide.]
77/11/03 Tom Kelly to Andy Mickel: [Burroughs Pascal from UCSD now running at Burroughs. 77/10/12 Tony Schaeffer to Andy Mickel: [Interactive I/O, language standards in the ligh of the natural evolution of Latin and ANS Fortran.]
77/08/25 Robert A. Fraley to Andy MickeI: [Comments on changing the definition of Pascal
77/11/07 Robert A. Fraley to Arthur Sale: [Comments on the Feat
conerning else in
77/12/26 Barry Smith to Andy Micke1: [Oregan Minicomputer Software history with ESI and
OMSI, and their PDP-11 Pascal implementation; Pascal T-shirts.]
77/12/12 Dave Thomas to Andy Mickel: [Pascal at Imperial College, London. A multi-user 77/11/07 Mitchell R. Joelson to Andy Mickel: [Law for the IBM 370 implementation at IC.] Mitall R. Joelson to Andy Mickel: [Law Enforcement Assistance Administration regulations vis-a-vis programming alnguages for use in criminal justice information systems.]
77/12/30 Ken Robinson to Andy Mickel: [Australasian distribution; Pascal use in Australian Universities, sundry comments on Pascal]

Special Topic: Pascal Standards: Introduction by Andy Mickel and Jim Miner: [ISO Standard Pascal, Conventionalized Extensions, Laundry Lists of Additional Features, Pascal Compatibility Report.]
77/12/09 Bengt Nordstrom to Andy Mickel: [The Swedish Technical Committee on Pascal;
Yet Another Attention List, will be in touch with the British Standards group.] Pascal Workshop with representation from industrial firms, governmental for and "academic experts". Consideration of DOD-1, a proposed Pascal-X extended version

Implementation Notes: General Information; Applications; Portable Pascals, Pascal-P4 Bug Reports and how Pascal-P4 relates to the standard. Pascal Variants: Pascal-S, Concurrent Pascal, Modula. Feature Implementation Notes: Unimplementable Features Warning; Compiling Boolean Expressions -- The Case for a "Boolean Operator Interpretation; Long Identifiers; Interim Report - Implementation of For Statement 2 More on For Statement. Machine-Dependent Implementations: Alpha Micro AM-100, CDC 6000 /Cyber 70,170 , CDC 7600 , Data General Nova/Eclipse, DEC PDP-8, DEC PDP-11, DECUS Pascal SIG; DEC LSI-11, DEC 10, HP-21MX, Honeywell 6000, level 66, H316, IBM $360 / 370$, ICL Clearing House, ICL 1900, ICL 2900, Intel 8080, MITS Altair 8800 , Motorola 6800, Prime P-300, Univac 1100, Zilog Z-80.

Pascal News $\# 112$, June, 1978, Pascal User's Group, University of Minnesota Computer Center 135 pages ( 70 numbered pages), edited by Andy Mickel.

Editor's Contribution: Personal Observations regarding Pascal-P, the first good critical article about Pascal, the need for a "business-oriented" Pascal procedure library, and more news needed about teaching experiences about Pascal. Status of Pascal User's Group: must raise rates for US and UK; rates lowered for Australasia.

Here and There With Pascal: Pascal Jobs, Tidbits (7 pages), French/English, English/French Pascal.Identifiers, Pascal in the News, Conferences, Books and Articles: Applications, Implementations, Languages, Textbooks, Reviews, Articles wanted. Roster Increment (78/04/22).

Applications (new section): News: Empirical study of Pascal programs (Pascal program style analyzer), numerical library project. Algorithms: A-玉 Random Number Generator A-2: Timelog; Software Tools: S-1 Compare (compare two textfiles for equality), S-2-1 Augment, S-2-2 Analyze (Pascal performance measurement programs); Programs: P-1 Printme (reproduce self).

Articles
"Extensions to Pascal for Separate Compilation"
Richard J. LeBlanc
[The lack of features in Pascal to allow procedures and functions to be compiled This weakness is particularly evident when modifications are being made only to limited parts of a program. Modificaitons of this sort are common, for example in the maintenance or extension of a Pascal compiler. By creating a global environment, separate compilation of routines using that environment, and additions to the environment without requiring recompilation of existing routines and declarations--all via extensions--a useful mechanism can be attained.]
"What Are Pascal's Design Goals?"
Robert Vavra
[As a long-time reader of Pascal News, the author has enjoyed the many articles in which people have discussed various features which could be added to Pascal, but they have been unable to take seriously. In arguing for or against some particular feature, writers have rarely involked Pascal's design goals in support of their
arguments. Such failure to build a proper foundation for one's arguments might be acceptable in casual conversation, but not in a serious discussion.]
"Pascal Environment Interface"
Terje Noodt
[Work is presented for a Pascal implementation for the Norsk Data Nord 10, running interactively. The Pascal Report does not say too much about how to interface a complier to a computer system and its users. To further complicate matters, what
it does say about this relates to a batch system, and is worthless or unusuable in an interactive system. A language is often judged on the way a particular
implementation interfaces to its environment such as what tools are available for the construction, compilation, and execution of a program, and what interfaces are like between the implementation and other systems on the computer (particularly the environment pollute Pascal!]
"Subranges and Conditional Loops"
Judy M. Bishop
[The subrange facility in Pascal is an aid to run-time security for fixed-boundary constructs such as counting (for) loops and array subscripts. The relevant types can be precisely and naturally defined, and the compiler can minimise the amount of run-time checking required. However, an index which increases under program control, as in a conditional (while) loop, presents a problem. This note
problem and presents a solution terms of a naming convention

## "A Few Proposed Deletions"

John Nagle
[Since quite a number of extensions to Pascal have been proposed, I thought that it would be desirable to propose a few deletions to keep the size of the language down With the goal in mind of keeping Pascal a simple, elegant, and useful language requiring a minimum of run-time machinery, I propose a few simple changes in the
direction of simplicity.]

Open Forum:
78/01/18 Arthur Sale to Andy Mickel: [Pascal News distribution in Australasia; explanation $77 / 11 / 11$ Giurge size of PUG(AUS) fee. Pascal as a first language in Australian uni's.] 78/02/02 Jerry reading and writing scalars, varying length strings, formatted input, ett.] 78/02/02 Jerry Pournelle to Andy Mickel: [Acquiring Pascal for a 48 K Z-80.]
18 last issue: doesn't miss comments to proposed extensions by Robert Fraley
78/02/23 Hellmut Weber to Andy Mickel: [Wish list from a usi's included files nice.]
ish list from a user's point of view for Pascal

78/02/24 Arthur Sale to PUG membership: [Commentary on Pascal News No. 11; David Barron's proposal for algorithms excellent, Pascal is not up for grabs, PUGN maturing.]
$78 / 02 / 27$ Greg Wetzel to Andy Mickel: [Shame on you for including Fraley's article-i
scared us--congratulations, you were terrifyingly successful! Stand by your guns.] 78/03/06 Eric Small to Andy Mickel: [Looking for Pascal programmer for consultants in broadcasting technology.]
78/03/08 Bob Jardine to Andy Mickel: [Reply to criticism of B6700 by Arthur Sale's Feature Implementation Note on Unimplementable Features.]
8/03/10 K. S. Bhaskar to Andy Mickel: [Pascal needs standardization and perhaps an
78/03/15 Terje Noodt to Andy Mickel: [A ne
interface is an important consideration.
78/03/16 Don Terwilliger to Andy Mickel: [Even though Tektronix is actively using Pascal it does not currently have products incorporating Pascal programming capabilities. 78/03/16 Edward Reid to Andy Mickel: [Interested in Arthur Sale's comments about Pascal 78 the B6700; comments on other items in past Pascal News issues.]
88/01/02 Werner Remmele to Andy Mickel: [Pascal implementation on the Intel 8080 using the ISIS II operating system. Notes about the project.]
78/03/15 Mark Horton to Andy Mickel: [Pascal at the Univ. of Wisconsin; comments about proposed extensions to Pascal, some more proposed extensions to Pascal
8/04/11 Jon Squire to Andy Mickel: [Pascal and DOD-1; need for a standard set of acceptance test programs for Pascal.]

7 Judy Bishop to John Strait: [A further comment on predefined types and subranges used in conditional loops.]

## ascal Standards

Introduction by Andy Mickel and Jim Miner: International Working Group by J $\phi$ rgen Steensgaard-Madsen investigating conventionalized extensions at last; News from Tony Addyman on the BSI/ISO Pascal Standard; criticism of the upcoming UCSD Workshop on Pascal Extensions. Naur Form on only 2 pages.]
8/02/06 Tony Addyman to Andy Mickel: [New phone number, urge that all PUG members commen on the BSI/ISO draft standard document.
8/02/01 Tony Addyman to DPS/13/4, Swedish Technical Committee and all correspondents:
8/03/2 [Update on progress by the BSI working group DPS/13/4 for a Pascal standards document proposed Morn Wher proposed summer Workshop at UCSD.
 format of Ken Bowles's summer Workshop on Pascal extensions.]
8/03/30 Bob Vavra to Andy Mickel: [Comment on Pascal's Design Goals; optimistic about the future of Pascal in spite of all the moves to extend Pascal.]

Implementation Notes: Checklist (new item); Portable Pascals (more Pascal-P4 bug reports) Feature Implementation Notes: Representation of Sets; Machine-Dependent ImplemenAutomation LSI-2,4, Data General Eclipse, DEC PDP-11, VAX 11/780, HP-2100, 21 MX , HP-3000, IBM $360 / 370$, Intel 8080 , Interdata $7 / 16,8 / 32$, Northwest $85 /$ P, Prime $\mathrm{P}=400$, Index to Implementations for issues 9-12

Roster Increment (78/10/31)

Following is a list of PUG members who either joined or changed address since the last roster increment was printed on 78/04/22. The list actually includes some persons who renewed, but whose address didn't change. Sorry.

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| 15238 | ROBERTA WACHTER/ INDUSTRY SYSTEMS DIVISION/ 200 BETA DRIVE/ PI TTSBURGH PA 15238/ (412) $782-1730 \times 544$ |
| 15260 | alan m. lesgold/ lrdc computer facility/ univ. of pittsburgh/ 39390 'hara st./ pittsburgh pa 15260/ (412) $624-4901$ |
| 16057 | peter richetta/ 287 Normal avenue/ SLIPPERY rock pa 16057/ (41 2 ) 794 -3531 |
| 16701 | FRANK BREKSTER/ 1 N. VISTA AVE./ BRADFORD PA 16701/ (814) 368-6319 |
| ${ }^{17331}$ | MICHAEL D. brown/ r. H. Sheppard co. inc./ 101 Philadelphia st reet hanover pa 17331/ (717) 637-3751 |
| 18016 | Charles t. Lewis/ bethlehem steel 1581 Martin tower/ bethlehe m pa 18016/ (215) 694-6359 |
| 18017 | ROBERT COLE/ 782 BARRYMORE LANE/ BETHLEHEM PA 18017/ (215) 865 -6509 |
| $\begin{aligned} & 18103 \\ & 18914 \end{aligned}$ | RICHARD J. CICHELLI/ 901 Whittier DRIVE/ ALLENTOWN PA 18103/ ( 215 ) 797-3153 (HOME)/ (215) 253-6155 (HORK) PHILIP W. ROSS/ 8 HICKORY LANE/ CHALFONT PA 18914 |
| 19004 | JIM Shallou/ 115 BIRCH AVE./ BALA CYYWYD PA 19004 |
| 19020 | ROBERT H. TODD JR/ Briarhood $11167 /$ CORNWALLIS HgT PA 19020/ ( 215 ) $752-4604$ |
| 19020 | BOB LIDRAL/ 3806 BENSALEM BLVD. $1214 /$ CORNWELLS HTS. PA 19020 |
| 19044 | JAMES A. MGGLINCHEY/ 296 BLAIR MILL RD. APT B-7/ HORSHAM PA 19044 |
| 19047 | RODNEY MEBANE/ 600 OLD STREET ROAD AAT13/ TREVOSE PA 19047 |
| 19083 | T. L. (FRANK) PAPPAS/ 338 Frascis drivel havertown pa 19083/ ( 215 ) $789-3206$ |
| $\begin{aligned} & 19102 \\ & 19111 \end{aligned}$ | RICHARD L. DAY/ TMME SHARE SUPPORT GROUP/ BELL TELEPHONE OF PE NNSYLVANLA/ ONE PARKWAY/ PHILADELPhIA PA 19102 ALAN M. KANISS/ 1327 MCKINLEY ST./ PHILADELPHIA PA 19111/(215) 441-2051 (HORK) |
| 19117 | dan morton/ 359 Northwood ave./ philadelphia pa 19117 |
| 19122 | BILL CHESWICK/ COMPUTER ACTIVITY - SYSTEMS GROUP/ TEMPLE UNLV. / BROAD \& LDNTGOMERY STREETS/ PHILADELPHIA PA 19122/ (215) 787-8527 (HORK) (215) 862-2153 (HOME) |
| 19128 | John F. RAATII 300 Hermitage st / Philadelphta Pa 19128 |
| 19145 | Paul J. pantano/ 2323 S . 17 ST ./ Phillidelphia pa 19145 |
| 19172 | mick cietkovic/ vic 6e/ penn mutual life/ 510 halnut st./ Phil adelphia pa 19172 |
| 19317 | BOB KELLER/ CONCORD WAY/ Chadds ford Pa 19317 |
| 19342 | JAMES I. hILLIAMS/ RD 4 BoX 18/ GLEN MILLS PA 19342/ (215) 648 -3554 |
| 19380 | THOMAS J. AHLBorn/ depr. Math/ hest chester state collegel hes t chester pa 19380/ (215) $436-2181$ |
| 19380 | GARY L. heigel/ 202 hestbrook drive/ hest chester pa 19380/ (2 15) 328-9100 (hork)/ (215) 696-8739 |
| ${ }_{19422}^{19401}$ | bILL BRENNAN/ 39 JODY DRIVE/ NORRISTOWN PA 19401/ (215) 277-24 66 <br> PETER A. NAYLOR/ MS b/220M/ SPERRY UNIVAC/ P.0. boX 500/ blue bell pa 19422/ (215) 542-3732 |
| 19422 | J. P. M. STOFBERG/ MS B/220M/ SPERRY UNIVAC/ P.0. BOX $500 / \mathrm{ELU}$ E BELL PA $19422 /$ (215) $542-4011$ |
| 19446 | michall rosiak/ 122 ardwick terrace/ Lansdale pa 19446 |
| 19446 |  |
| ${ }_{19711}^{19518}$ |  |
| 19711 |  |
| 19898 | Samuel c. kahn/ info systras dept n-1450 planning div. / du po nt co./ hilmington de 19898 |
| 20005 | John b. holmblad/ TELENET COMMUNICATIONS Corp./ 1012 14TH ST. NH/ hashington dc 20005/(202) 637-7900 |
| 20016 | RICHARD B. FITIL/ 4215 38TH STREEE NH/ WASHINGTON DC 20016 |
| ${ }_{20037}^{20022}$ |  |
| 20052 | MICHAEL B. FELDMAN/ DEPT. OF EE \& CS/ GEORGE WAShington univ./ hashincton dc 20052/ (202) 676-7593 |
| 20052 | e. michael hamilton/ C.A.A.C./ george hashington univ./ 2013 G Street nh/ hashington dc 20052/ (202) $676-6140$ (202) $676-3673$ |
| 20250 | T. Q. STEVENSON/ O \& F DATA SERVICES/ RM 4646-S/ USDA/ WASHING TON DC 20250/ (202) 447-6275 |
| 20375 | NIELS X. WINSOR/ CODE 6752/ NAVAL RESEARCH LABORATORY/ WASHING TON DC 20375/ (202) 767 -3134 |
| 20755 |  |
| 20810 | RANDY BARTH/ 9206 CANTERBURY RIDING/ LAUREL MD 20810 |
| 20822 | tom entrrime/ 13311 Chauscey pl. f203/ mi. rainier md 20822 |
| 20850 | TOM LOVE/ SOFTHARE METHODOLOGY/ GENERAL ELECTRTC/ 401 N. HASHI NGTON ST./ ROCKVILLE YD 20850/ (301) 340-4000 |
| 20852 | ATTN: INFormatics Inc. booxstore/ 6011 EXECUTIVE BLVD./ Rockvil Le in 20852 |
| 20853 | PATRICIA SHELLIL INFORMATICS |
| 20855 | Bob Rogers/ 18625 AZALEA DRIVE/ DERHOOD MD 20855/ (301) 869-20 89 |
| 20903 | H. A. Cook/ 1223 CRESTTAVEN DR./ SILVER SPRING MD 20903 |
| 20904 21030 | John g. guthriel computer entry systens Inc./ 2141 INDUSTRIAL PARKHAY/ SILVER SPRINGS Md 20904/ (301) 622-3500 |
| ${ }_{21043}^{21030}$ |  |
| 21045 | barton f. Norton/ Chromal p.o. box 126/ CoLUMBIA MD $21045 /$ (30 1) $992-7404$ |
| 21203 | ROB BIDDLECOMB/ MS 451/ hestinghouse Electric corp./ sdd east box 746/ baitimore mi 21203/ (301) 765-6322 |
| 21204 |  |
| 21235 | LeSter sachs ms 3 -0-25 OPER. blde/ Social security administra tion/ 6401 SECURITY boulevard/ baltimore md 21235/ (301) 594-5482 |
| ${ }_{21701}^{21401}$ |  |
| 22030 | ATTN: J. . . P. Associates/ 3219 Prince hillimm dr./ Fairfax va $22030 /(703) 591-8525$ |
| 22101 | david auli/ Cotputer science/ wp 615/ The mitre corp./ 1820 do LLY Madison blvd./ Mclean va 22101/ (703) 437-7898 (home) |
| 22101 | H. F. HESSION/ ADVANCED RECORD SYSTEMS ENGINEERING/ KESTERN UN ION/ 7916 HEST PARK DRIVE/ MCLEAN VA 22101/ (703) 790-2241 |
| ${ }_{22151}^{2215}$ | ROBERT A. GIBSoN/ 8902 NTCOL LANE \%207/ MANASSAS VA 22110/ (70 3) 367-4792 (HORK)/ (703) 369-5640 (HOME) |
| ${ }_{22151}^{22151}$ | harold d. Jenkins Jr./ Springeield support center/ fairfax cou nty public schools/ 6707 electronic dr./ springeield va 22151 PALL T. DYKE/ RESOURCE SYSTEM \& PROGRAM ANALYSIS/ 428 GHI BLDG ./ U.S.D.A./ 500 12TH ST S.h./ WASHINGTON DC 22151 |
| 22180 | Robert g. Fitzeeraid/ 133 East street n.e./ vienna va $22180 /$ ( 301 ) $868-5229$ |
| 22205 | halter a. hhite/ 6048 N 9th st./ ARLTNGTon va 22205 |
| 22209 | LARRY DUBY/ 1500 ARLINGTON BLVD. finlo/ ArLITGTON VA 22209 |
| 22309 22310 | GERALD P. SHABE/ 3206 NoRhtch TERRACE/ ALEXANDRIA VA 22309/(703) $360-5587$ |
| 22310 22312 |  |
| 22312 22401 |  |
| 22801 | MICHAEL STAUFFER/ EASTERN MENYONTTE COLLEGE/ HARRISONBURG VA 22801 |
| ${ }_{22923}^{22901}$ | AVERY CATLIN/ THIMBLE FARM/ ROUTE 5 - BOX 363/ CHARLOTTESVIL* VA 22901 <br> LINHOOD FERGUSON/ RT 1 bOX 3C - LAXR SAPONI/ BARBOURSVILLE VA 22923/ (804) 973-516 |

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PASCALNEWS #13
DECEMBER, 1973
\begin{tabular}{|c|c|}
\hline 23185 & glas dunlop/ 1502 convay drive - apt. 103/ williamsburg va 23185/ (804) 826-1 \\
\hline 23284 & Agnes h. Elmore/ Computing activities/ virginia commonwelith unvy / 1015 floyd ave./ richmond va 23284 \\
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\hline 23505 & R. E. CRItTSIMGER JR./ 136 BLAKE ROAD/ NORFOLK VA 23505 \\
\hline 23505 & LLOYD D. FINK/ ATR CARGO INC./ P.0. Box 9793/ NoRFoLK YA 23505 / (804) 480-2660 \\
\hline 23669 &  \\
\hline 27702 & hilliam h. diugiud/ planning div./ city of durham/ 101 ctit ha ll plaza/ durbam nc 27702 \\
\hline 28214 & warren c. Fordham/ Mcclure lumber co./ 6000 MT. Holly rd/ Char lotte nc 28214 \\
\hline 28704 & Carroll b. Robbins Jr./ Apt 32/ ARDEN ARMS APTS./ Arden nc \(28704 /\) (919) 684-0168 \\
\hline 29206 & HOWARD EISENSTEIN/ 6616 DARE CIRCLE/ COLUMBIA SC \(29206 /\) (803) 782 -5041 \\
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\hline 30021 & Craic m. Inclis/ 1420-C post oak dr./ clarkston ga 30 \\
\hline 30033 & John p. Cuches/ The hyde company/ 2169 CLATPMont rd ind decatu r ga 30033 \\
\hline 30067 & HENRY D. KERR III/ 4820 HAMPTON LAKE DRIVE/ MARIETTA GA 30067/ (404) 971-2197 \\
\hline 30303 & darreli preble/ computer center user services/ georgia state u niversity/ atlanta ga 30303/ (404) 658-2683 \\
\hline 30305 & JEFFREY H. BIGGERS/ SUITE 411/ DTh INC./ 3100 MAPLE DRIVE NE/ ATLANTA GA 30305 \\
\hline 30305 & hilliam g. christian/ sutte 450/ CLS inc. 3100 maple drive ne / atlanta ga 30305 \\
\hline 30305 & Frank s. SPARKMAN/ SUTTE 411/ DTW Inc./ 3100 Maple drive ne/ a tlanta ga 30305 \\
\hline 30305 & DAVID T. WILSON/ SUITE 411/ DTH INC./ 3100 MAPLE DRIVE NE/ ATL ANTA GA 30305 \\
\hline 30327 &  \\
\hline \({ }_{30354}^{3034}\) &  \\
\hline 32204 & ATtention: roy h. Fileger/ sutte 110 EAST/ Computer power/ 661 RIVERSIDE AVE/ Jacksonville fl 32204 \\
\hline 32304 & prgGy roblyen/ Eddcational computing project/ florida state de pr. of education/ tallahassee fl 32304 \\
\hline 32407 & ANNA WATSON/ 3705 delwood drive/ Panama city fl 32407/ (904) \(234-4423\) \\
\hline 32670 & RTCHARD J. NAST/ 1721 SW 55Th Lane/ ocala fl 32670 \\
\hline 33065 & HOWARD S. MARSHALL JR./ 2648 NW 86TH AVE./ Coral springs fl 33065 \\
\hline 33068 & dean james/ 7440 S.W. 10TH ST. - \(1102 / \mathrm{N}\). LAUDERDALE Fl 33068 \\
\hline 33142 & MONTE ELLIS/ RYDACOM INC./ 3401 NV 36TH ST./ MIAMII FL 33142 \\
\hline \({ }^{33181}\) & JAMES GROSSMAN/ 2365 MAGNOLTA DR./ N. MIAMT FL 33181/ (305) 89 1-3440 \\
\hline 33528 & Clara l. johnson/ media research div. - engineering a. C. nie lson co./ 375 Patricia ave/ dunedin fl 33528/ (813) \(734-5473\) \\
\hline 33601 &  \\
\hline 33803 & ALLEN F. DOWNARD/ 3008 REDHOOD AVE./ Lakeland fl 33803 \\
\hline \({ }^{35801}\) & Marvin e. kirtti/ 1327 monte sano blvi. S.E./ huntsvilue al \(35801 /\) (205) \(837-7610\) \\
\hline 35803 & DAVID MCQUEEN/ 2410 ARROYLDOD DR./ HUNTSVILLL AL \(35803 /(205) 881-3628\) \\
\hline 37076 & LARRY D. BoLES/ 649 DENVER DRIVE/ Herkitage TN 37076 \\
\hline 37660 & J. K. DISSELKAMP/ 202 building 54/ TENNESSEE EASTMAN COMPANY/ kIngsport tn \(37660 /(615) 246-2111\) \\
\hline 40206
40583 &  \\
\hline 43147 &  \\
\hline 43229 & RICHARD E. ADAMS/ 967 ATLANTIC AVE \(7634 / \mathrm{COLUMBUS}\) OH \(43229 /\) ( 6 17) 436-3206 \\
\hline 43230 & RICHARD GREENLAS/ 251 CoLony court/ gahande oh 43230/ (614) 475 50172 \\
\hline 43762 &  \\
\hline 43778 & toun legrazie/ rural route 1/ Salesville of 43778 \\
\hline 44092 & LYNN C. HUTCHINSON/ BALLEY Control company/ 29801 buclid ave/ wickitfee oh 44092/ (216) 943-5500 \\
\hline \[
\begin{aligned}
& 44106 \\
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\] & JACK D. ALANEN/ JENNINGS COMPUTING CENTER/ CASE HESTERN RESERV E UNIV./ Clevelind of 44106/ (216) \(368-2800\) \\
\hline 44106 & PAUL MEILAND/ DENTAL SCHOOL CLINICS/ CASE WESTERN RESERVE UNIV .12123 ABINGTON ROAD/ CLEVELAND oh 44106 \\
\hline 44107 & BILL SHANNON/ 2038 arthur/ LakEhood oh 44107 \\
\hline 44512 & ATTN: hestern reserve communications/ 424 INDIANOLA ROAD/ Young stown oh 44512 \\
\hline 45201 & hilliam r. metz/ MSD - develophent The procter \& gamble compa ny/ f.o. box \(599 /\) cincinnati oh 45201/ (513) 562-2747 \\
\hline 45215 & G. D. Montillon/ 351 Flekmigrd./ cincinnati of 45215 \\
\hline 45241 & francis h. bearden/ data sxsteks/ cincinnati electronics corp. / 2630 glendale-Milford road/ cinctnnati oh 45241/ (513) 563-6000 x140 \\
\hline 45244 & CLINTON HERLEE/ MEDIATOR INC./ 2812 SADDLEBACK DRIVE/ CINCTNNA TI OH 45244 \\
\hline 45342 & D. R. HILLL MONSANTO RESEARCH CENTER/ P.O. Box \(32 /\) MTAMMSBURG OH 45342 \\
\hline 45409 & DAA C. WATSON/ WRIGHT BROS./ BOX 541/ DAYton oh 45409/(513) 2 23-2348 \\
\hline 45414 & Lahrence a. Shively/ 6014 Frederick roadj dayton oh 45414 \\
\hline 45424 & M. B. CLAUSING/ 5603 FISHER DRIVE/ DAYYON oh \(45424 /\) / (614) \(236-3475\) \\
\hline 45424 & W. A. Shull 4063 BUTTERROOD Court/ Dafron on \(45424 /\) (614) \(233-6487\) \\
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46201 & JoE CLMA/ SUITE 200/ SIMULATION TECHNOLOGY INC./ 4124 LINDEN AVE./ DAYTON OH \(45432 /\) ( 513 ) \(252-5623\) \\
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\hline 46205 & RICHARD A. byers/ 3690 GLENCAIRN LANE/ INDIANAPOLIS IN 46205 \\
\hline 46312 & vincent elias/ Security federal s \& l assk./ 4518 Indinapolis blvd./ east chicago in 46312 \\
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\hline 46526 &  \\
\hline 46805 & R. GARY LeE/ DEpr. OF COMPUTER TECHNOLOGY/ PURDUE UNIV./ 2101 CoLISEIM BLUD./ Fort hayse in 46805 \\
\hline 46808 & DALE GAMMER/ GOVT. \& INDUSTRIAL DIV./ MAGNAVOX/ 1313 PRODUCTIO N ROAD/ FORT HAYNE IN 46808/ (219) 482-4411 \\
\hline 47272 & Donald l. Clapp/ R. \$1/ ST. PAuL IN 47272 \\
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\hline 47907 & Edhard f. gehringer/ dept. of computer science/ math sciences building/ purdue universtiy/ h. Lafayette in 47907/ (317) 743-3 \\
\hline 47907 & SAUL ROSEN/ Computing center/ G175 math sciences bldg/ purdue univ. h. Lafayetre in 47907/ (317) 494-8235 \\
\hline 47907 & MIChaEl deisemroth/ School of ind encr. / Purdee univ./ h.Lafay ette in 47907/ (317) 493-3157 \\
\hline 48010 & Shaun devirin 6854 CEDARRroox/ birkt ingham mi 48010 / (313) 322-685 \\
\hline 48033 & H. DICK Breidenbach/ 4955 PATRICK/ W. BLOOMFIELD MI 48033 \\
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48098 & CHRISTOPHER A. PHILLIPS/ 29205 LUND / SOUTH BLDG. APT 14/ WARR EN MI 48093 WESLEY E. MANGUS/ 5786 NORTHFIELD PKWY./ TROY MI 48098 \\
\hline 48103 & ALAN A. Kortesoja/ 701 h . dAvis/ and Arbor mi 48103/ (313) \(995-7063\) \\
\hline 48103 &  \\
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48103 & wILLIAM LUITJE/ 2509 WEST LIBERTY ROAD/ ANN ARBOR MI 48103/ (3 13) 769-7820 LES WARNER/ 1804 LINWOOD/ANN ARBOR MI 48103 \\
\hline 4818105 &  \\
\hline 48105 & KURT METZGER/ 478 Cloverdaie/ ann arbor mi \(48105 /\) (313) \(662-4757\) \\
\hline 48106 &  \\
\hline 48109
48130 &  \\
\hline 48169 & Join s. gourlay/ 8645 toma road Pinckney Mi \(48169 /\) (313) 994-6645 \\
\hline 48184 & DAVID R. POSH/ DEPT \(3741 /\) BURROUGHS CORP./ 3737 s . VENOY RD./ HAYNE MI 48184/ (313) \(722-8460 \mathrm{X} 267\) \\
\hline 48228 &  \\
\hline 48640
48824 & bob metzger/ Computer technology dev./ dow chemical co./ 2040 dow center/ midland mi 48640 J. F. p. MARCHAND/ CYCLOTRON LABORATORY/ MICHIGAN STATE UNIV./ EAST LANSING MI 48824 \\
\hline \({ }_{49001}^{4824}\) &  \\
\hline 49003 & James h. haiters/ American national holding co. P. P.O. box 2769 / kaiamazoo mi 49003/ (616) 383-6700 \\
\hline 49006 & LOREN L. HEUN/ THE UPJOHN CO./ 301 HENRIETTA ST. / KALAMAZOO MI \(49006 /\) (616) \(385-7886\) \\
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49008 &  \\
\hline 49085 & ANTHONY J. SCHAEFPER/ 1023 VINELAND RD./ ST. Joseph Mi 49085 / (616) 429-8517 \\
\hline 49464 & MIKE HAMMAN/ HERMAN MILLER INC./ ZEEMAN MI 49464/ (616) 772-33 00 \\
\hline 49503 &  \\
\hline 49931 & kenneth m. mCmilin/ simulation lab/ mchican tech univ./ hovg hoon mi 49931/ (906) 487-2111 \\
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\hline 50158 &  \\
\hline 50307
51106 &  \\
\hline 52240 &  \\
\hline 52302 & don stover/ 2270 26TH Street/ MARIoN IA 52302/ (319) \(377-8529\) \\
\hline 52302 & dennis sutherland/ 283525 TH AVE./ MARION IA 52302/ (319) 377-4200 \\
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\hline \({ }_{5}^{53201}\) &  \\
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\hline \({ }_{53218}^{53217}\) &  \\
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53927 & O. ARTHUR STIENNON/ PARK-REGENT MEDICAL BLDG./ 1 SOUTH PARK ST ./ KADISON WI 53715/ (608) 255-6262 JAMES E. TARYID/ BOX 20/ DELLHOOD WI 53927/ (608) 339-7259 \\
\hline \[
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& 54601 \\
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\] & \begin{tabular}{l}
THOMAS C. HICKS/ 1108 S 5TH. ST./ LA CROSSE WI 54601/ (608) 78 4-4345 \\
AND ÉK L. PERRIE/ 1208 BAY SHORE DR./ OSHKOSH WI 54901/ (414) 233-4661 (HOME)/ (414) 424-2068 (HORK)
\end{tabular} \\
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DAVID P. MCDONNELL/ 4912 GEDDES/ FT. WORTH TX 76107/ (817) $738-1884$
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DAVID HOLLAND/ SUITE 200/ INTERCOMP/ 1201 DAIRY ASHPORD/ HOUST ON TX 77079
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(
GENE HUGHES/2907 SKY CLIFF/ SAN ANTONTO TX $78231 /$ (512) 492-9 661
ATTN: COMPUTATION CENTER/ TEXAS A \& I UNIVERSITY/ CAMPUS BOX 1 85/ KINGSVILLE TX 78363
KATIE NOONING/ 1105 DEEPWOOD/ ROUND ROCX TX 78664 ( 512 ) $255-6$ O 052 (
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ALAN ZARING/ COMPUTER SCI. DEPT./ UNIV. OF TEXAS - AUSTIN/ AUS TIN TX $78712 /$ (512) $471-7316$
DONALD G. WEISS/ H2565/ 3501 ED BLUESTEIN BLVD./ AUSTIN TX 787 21/ (512) 928-6034
ROGER H. FRECH/ 10033 CHILDRESS DR./ AUSTIN TX 78753/ (512) 83 7-6078
WALT FEESER / 8900 SHOAL CREEK SUITE $109 /$ AUSTIN TX 78758
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HARRY P. HAIDUK/ 6202 MCCOY/ AMARILLO TX $79109 /$ (806) (915-5111
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thomas l. LIGHT/ STORAGE TECHNOLOGY CORP./ 2270 S 88 TH ST./ LO UISVILLE CO 80027/ (303) 666-6581
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bill ehlert/ p.o. box 3154/ DENVER CO 80201
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PASCAL NEWS #13
DECEMBER, 1978
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PASCAL NEWS \#13
DECEMBER, 1978
PAGE 28


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| Roberct l. Jardine | 92691 | HESLEY E. MANGUS | 48098 |
| george d. jelatis | 55417 | J. F. P. MARCBAND | 48824 |
| HAROLD D. JENKINS JR. | 22151 | thomas a. marciniak | 20853 |
| JOHN JENKINSON GREGORY JENNINGS | 75006 53207 | ${ }_{\text {RICK L }}$ L. MRES M MARCUS | 55455 |
| Jefrrey c. Jenning | 59801 | GAYE MARR | 01867 |
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|  | ${ }_{7} 90274$ | James a. mgelinchey | 19044 |
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| michael h. leskin | 10021 | THOMAS M. NEAL | 92634 |
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| Ping K. LiAO | 94545 | ROBERT C. MICKERSON | 94611 |
| KARL LIEberbirr | 87135 | StuAR C. WIMS | 90266 |


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| barbara K . nukth | 13760 |
| BARTON F. Morton | 21045 |
|  | 60626 |
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| htlliam J. nyback | 60606 |
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| robert Osborm | 02178 |
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| JAMES K . $0^{\prime}$ Connor | 07332 |
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| G. 0'schenectany | 12202 |
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| dennis paull | 94022 |
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| eric peabody | 75042 |
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| ${ }_{\text {Hat Perso }}$ | 55102 |
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| h. hesley peterson | 96822 |
| Christopher a. philurps | 48093 |
| PAUL PICKELMANN | 48109 |
| DAVID PICKENS | 80302 |
| WILL PICKLLES | ALI3 4RZ Unitimd krngdom |
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| halter l. pragnell | 02149 |
| darreli preble | 30303 |
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| RICHARD E. PRICE | 60630 |
| Fhllimm c. price | 97068 |
| hilison t. price | 94619 |
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| John L. PRUN | 90630 |
| EPIIC PUGH | 90024 |
| charles j. purcell | 55113 |
| James l. PYuEs | 02181 |
| steve qualline | 13210 |
| rapmond quiring | 90019 |
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| M. Rahilly | 3072 AUSTRAL |
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| ROBERT J. reymolds | 92138 |
| d. Lloyd rice | 90406 |
| peter richetta | 16057 |
| dAN RICHMOND | 92103 |
| george h. richmond | 80027 |
| Lorin ricker | 97225 |
| Charles rider | 91326 |
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| arroll b. robbins jr. | 28704 |
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| JoE c. roberts | 75042 |
| TERRY R. Roberts | 80202 |
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| HOHARD RUMSEY JR. | 91105 |
| $\underset{\text { Martin ruwyan }}{\text { LeSTR }}$ | ${ }_{21235}^{60120}$ |
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| DECEMBER, 1973 |  |  | PAGE 29 |
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| inving s. schechtman | 08540 | KENNETH R. HADLAAND | 01420 |
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| Jahes r. Sccirage | 08854 | R. L. hallace | 92675 |
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| robrrt schutz | 11756 | R. WALHYN | 22 9NF UNITED KINGDOM |
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| JaCk p. Shal | 53012 | PER-AKE WEDIN | 90187 SkEben |
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| mike tiller | 55116 |  |  |
| ${ }^{\text {RoN }}$ TrPTon | 54134 |  |  |
| catherine c. tobey | 91335 |  |  |
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| p. Torgrimson | 94087 |  |  |
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## Applications

Please send all contributions for this section to Rich at the address below.
software tools


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ANPA/RI and
Lehioh University

THE MLONG AGO" PAST

In PN\#6 of November 1976, I introduoed the idea of ounversal pasen col Softwore Tools set, Bequuse the software Tools section is now a part
of pN, ond pug membership has ingreased by a faetor of five since No: vember 1976, it is relovent to remstate the ideas developed then,

SOFTWARE TOOLS FOR PASCAL
(*Fram "Pascal Potpouri" Pascal Nowsletter \#6 Novembar 1976*)
Pascal implementations for new environmenta are occuprina With ever increasing frequengy. As paseal is used for more and more production prognamming, it la important that a universal bet of anclliary soffware tools be agreed upon, some of these tools can be defined in an enviponmenteindependent woy so ehot When writen in standard eascal, thoy oan bocome part of o une iversal pascal software dovelopment factily, the here propose
an initial list, Wish PUG momberahip holp, the list will dea velop into. Working opecificetion and epowerpul set of prem gramming toole.

PASCAL COMPILERS

Currently there exist pascal aqmpliens whiah produge absom lute code, relocotable coder macro code (Pascal=J) and inter. preted code (Rascolap) Portable versions extst (Pasoalmp and
Pascalaj). Compilor trunks exist, A standord pasoal subset Pascal-J) Compi
(Pascalas) extstes,

For compllar writers there should be atandard pacal language test sat. This universal set of pascal pragnama would tow pasal compllers and help implementors galn confle dence in the correctines of their oompllers,

An intoractive interpreter should be developed, this sys. tem would provide interactive symbolic run lime debuoging fea cilities: breakpoints, interactive dumps, ite, It should be easy to do better than plitita Cheokout qompilif.

## The Legonme and Bochmann compller writing systeme ore leo mportant tool for any shop engaged in lenguage development.

## SOURCE PROGRAM TOOL

Wirth has writton e eroas reterence program, pophops, if the varlable names wepe improved, etanderd vereton of this prognam could be among the aftwore toolb. A formater or pretty printer" is esential for produoing documentation qual

A oode inatrumenter is every impontant debugelna and ram ining tool. Instrumenters inserf atatement countors or timer -a that reporpe of pelative uage of gode oen be made, An ine etrumentor is invaluabie in opimizing programe,

Ahigh level macro preprocessor wouldelso be e velugble focllity.

## SOURCE LIBRARIES

The CDC source librony utility proorom UPDATE is ourrontly used tar diatribution of the SCOPE version of Pescal. It seems to me that a mintuversion of UPDATE (With only sequential prooram libraries) cauld be implemented lí pasoel. This would
 have ever seen. 1 think ita quality should be emulited,)

Fon truly laroe systoms (50,000t lines) a source code data base is desinable, such a syitem keeps track of which program access what data and provides for standard itlo and rocord dosm apiptions amang proopams, otc, I yndepstand aych a satom for
documentation preparation
W. Burger implemented part of Waltef PLAP In Pascal. We need a univenaal PLAPmilke tool to mainteln manuals and other dooumentotion in machine readable form, Justificotion and hyphenation and foolitites for producino hiah quallay printing In upper and lower case should exist, pascal documentation ifation and diatribution.

## OBJECT PROGRAM FACILITIES

Work is now in progress on programe which load pasoal abe solute binaries, Facilitios for overlay processing should be structures should be provided. A binary decoder salso usiem ful tool.

OTHER PROGRAMS
An efflofent table processor with fachlities like cobols Report Writer would be deeinable, current work on Pascal date oomputer idod instruction systeme aunction the day of increased use of poscol in business, engineering, end eduopiton. In the area of function librapies (for mathemetics of butiness), fom allities should be provided for not only linkino in binapy món dule: but also for including aurge moduleit*

CONCLUSIONS

Qbviously, where onyironmentel conditions permit we should hove - Univerial rasacal progrem Implementing eaeh software ald Where the onvironmental iactops prevent this, we should seek to provide a standard waer intertace to the desifed functions,
*In my opinian, meroing proppame at the source level ts to be proferred to blary levol linking, pascal compllers are typlcally fastor than linkingoloadens.

SOFTWARE TOOLS CONTP

I bellove that early article presented viable porspective for pur ture pascal Software Tools (PST) Work, but it left unanswered many impore tant polloy questions. The moat erlitioal of these was how to got PST to Pug members one obvious onswor was that implementors could dintribute the tools with their distribution pagkage, of oourse, when the number of Whet to do. About this time CACM stopped its Aloorithms soction unaure of most concelled my membepship because of my feelings that the most impore tant wopk of compuser scientists is their prognama. Without published programs, CACM is hollow for me,
 oumenţal
() Publishing quality pascal programe would help educate new PUB members, I belfove poding good prognoms is the oasiost way to leann prooramming teohniques and style.
2) Publishing programs would give proper reçagition to program au* thore.

4) Published PSTes would encourage implementora ta adhepe to the pascal standard.
5) Commercial ysers could require çompller vendors to use the pStes to sent the oonformity end performence of theip implementetions,
Andy wes convinced and y oven impressed myself with these aroumente The obvious result is the Software Tools section started in PN\#I2.

The above disoussion should help PUG members understand and shape PN
Applications

Of oourse, publishing programs has its own probleme. The prelimina ry statement of what we are trying to do was in PNil2i

We decided to aresto new section for ppinting pasoal source programe for vorlous applications including software of sionifigant appifications pragram wition in pasacal. Jim Miner suggested we index each program so that they may be easin ly referenced for corpection and giticisms,
ion. Arthur solo la vory enthushastic about the Algorithms sece Arthur salo la vory anthula for

1) The provision for oeftiflcation of the program by unrea lated pensons, with eleer dentifioation of the system used; and
2) Critiques of the program forl
a) standards conformanoe,
b) style,
c) algonithm
d) output convenience and general design.

Weill number prognams stapting with pai, software Toole
starting with $\operatorname{sml}$, and Algorithms starting with $A-1$,
Already our numbering schome is giving us problems. As Sec (Augmen and Analyze) made clear, software tool may not be Just aingle programa but entire systems of programs. so that wo con retor to text lines with In programe uniquely, we will ueo the notationit

Referencenumber = Closeificotion Men syetem


Thus, Augment and Analyze are still saz but Augment, feself, beaomea S-2-1 and Analyze s-2-2,

The validation sulte designators (l, e, "YM programs) are being ase igned by Bpion Wichmann, so far there have beon meny favarable comment haye beon received to date. It is important that those mapers who bping up the programs oomment on eny problems they might have had, We really need to know the performance and ease of installation of the programs on various systems.

I'm sure the membership would be vory interested in how well Jim
Miner's Compare program (Sel) poriorms ogoinste pascal implementatation Miner's Compare Program (S-i) porforms ogolnat e Poscal implementatation
 like a good term prolect to me,

## THE PRESENT

This issuers pretty printers should help futura pST submitters to produoe cameramepady copy, Wo are oxpermanting with several publication atylea, We want to be able to publish roadable copy of large programa

We may expeniment with puteing multiple simple stetemente per line and systome of an many as io, q日ø lines of code in ingole lisue, thui, full compllers and large appliaations libparles oould be occommodated, Anyone eaking to have large prognampublished should work closely with me and Andy. Ingidentally, programa which ore primopliy tutorial in nature (i,ep not of general utility) hould be ingorporated in articles for

## the immediate future

I'm suro it ia of no suppriae to any roader that elmost all of the softwone tools dezeribed in PN\#G are now in existence, Unfortunately or publication. We are orking on getting Ior publicilong we ere working on getting those to whioh wo have acgess thons and onganizations, Work is golng on in ell these enese

We belleve the software rool get will grow in two waysi now utilie les will be odded and extating utilities will be modified or roplaced by mproved versions, he onoourage the momborship to holp us caretully evan Wate published programs, We hepe all those who have developed pasea oftware Tools will try to ubmt them aper publicetion

## WHAT'S NEW?

Condict: pretty prinker allows us to publith program text in a aonm istont atyle, Equally importont for program ohoring ore a ource dode forary facflity end a toxt formatter for documontation. After these loh prioplty foms have been publizhod, PN will dazzle the PUG membere hip with extraordinary eftwere

In eddition to the sortware mentiened in the PNKG artiole, we hope so publish on APL interpreter wiltton pntiroly in Pascal ot vilionava, A

 program from North. Amerioan philliks copporation whigh reads pasol prow
arams ond mapks them for standord conformance. Tho program handios most,
 this program into peoogizing the tuli standerd should contoet me immedlo otely.

## WHAT ARE WE LOOKING POR?

An interactive editor in pascal would be most weleoma, A typenete ting packege would halp newaletter produckion immensoly. We've writtien Motonole 6000 Aseambler in pegcel at ANPA/RI, It and ones for the 8080,


I'd like to eee bootetrapoble verelon of pascales published in PN, Host programmers Would be iunprised ot Juet how'eagy it is to qompli: aseal, Reading nice pacel complior writan in Raseal would make very prognommer reticent to muck with the andend, ingldentelly, the

Lare Moseberg of Volvo fiygmotor in sweden pointed out the impore



## GOME FINAL WORDS

Those implementers and organizotions, whioh insist on producin heir own dialects of pascal and folsting them on on unsuspecting pubile as beino pascal, are encition of us all.

## ALGORITHMS

A-1 Random Number Generator (continued discussion)
University of Lancaster

> Department of Computer Studies
> Bailrigg, Lancaster
> Telephone Lancaster 6S20I (STD 0S24)

Head of Dcpartment:J. A. Llewellyn B.Sc., M.Phiil., F.B.C.S., F.I.M.A
7 th September 1978

Dear Rich,
Jim Miner made a few comments on my random number generator algorithm (PN $* 12$, algorithm A - 1) which I feel compelled to enlarge upon:
(a) I don't know what Jim meant by the results "seeming" better with circular left shift. The original algorithm has only one absorbing state (i.e. a state which you can't get out of once you are in it), which i the zero state, and this is isolated (i.e. the only way you can get algorithm has two absorbing states, one at least of which is nonisolated. In practical terms this means that unless you are very care ful about choosing your initial seed, you wind up repeating the same number.
(b) As long as overflow checking is suppressed, multiply overflow can be ignored. For, if the initial seed is positive, then $a, b$, acomp and bcomp are also positive; hence a' (after the first shift) is positive; thus acomp' is positive, and the result of ( $a$ ' and bcomp') or ( $b$ ' and acomp')
must always be positive, independent of the sign of $b$
(c) I take the point about set operations expressing exclusive-or's more naturally, though this is exploiting a feature available in that particular implementation - in our implementation, integers occupy one the system.
(d) I pointed out in my note iv), that the initial seed must be positive and non-zero.

I hope the above comments are sufficient to prevent anyone using the modified algorithm before its properties have been more fully investigated.

## A - 3 Determine Real Number Environment

DOCUMENTATION : ENQUIRY

Language : Pascal
Written : A.H.J. Sale
Monday, 1978 March 20

## Use

To allow programs to enquire into their environment (compiler + computer) and tailor their behaviour to the properties of the real arithmetic system The procedure may be of use in programs that must be portable across many different PASCAL systems, and which are numerically oriented

## User documentation

Calling the enquiry procedure with the proper actual parameters determines he base and number of digits of the mantissa of the representation, and an edure works on a large range of computers, its correct operation depends on a number of assumptions about the representation of real numbers, and the operation of floating-point arithmetic. Programmers incorporating the procedure into programs are advised to cause the deductions to be printed so that end-users can check the accuracy of the deductions for their particular systems.

## nstallation

The enquiry procedure is standard PASCAL, in reference language form, and should compile on all systems. If assumption (b) (iv) is violated (as for example on the IBM 1130 which has more mantissa digits in its software accumulator than in the memory representation), rewriting the parenthesized expressons and iorrect deductions for the memory representation. The same trick may be employed in defense against over-clever optimizing compilers that utilize properties of (mathematician's) real arithmetic, and re-organize expressions.
The displayed driver program illustrates how the best- and worst-case precision may be computed from the deductions about the arithmetic.

## System documentation

The algorithm is an adaptation of one originally due to M.A. Malcolm (Comm ACM, Vol. 15 No. 11 pp 949-951, November, 1972).

Assumptions
$t$ is assumed that
(a) Real numbers are represented by floating-point representations which comply with the following conditions.
(i). There is a mantissa of a fixed number of digits to a fixed base
(ii) There is an exponent which expresses a multiplying factor to be pplied to the mantissa to obtain the exact representation multiplying factor is the base to the power given by the exponent.
(iii) The representation preserves maximum precision (no digits are lost unless the representation cannot accommodate them). In lost unless the representac with possible exact representations are exactly represented.
(b) Real arithmetic complies with the following rules
(i) If operands and results are exactly represented integral values no inaccuracy is introduced by the arithmetic.
ii) The arithmetic is organized along the usual align, operate and normalize steps, where these are necessary.
(iii) It is presumed that when digits are lost due to the represent ation, they are either truncated (ignored), or true rounding ation, they are either truncated (ignored), or true rounding
takes place. (No other possibilities are taken into consider ation.)
(iv) The intermediate results of arithmetic operations are held in a cell which has the same representational properties as the operands.
 SGET S LIST LINEINFO STANDARD
Programinvestigaterepresentation; var
> base,
numberofdigits
iounding
round

intoger;
$\vdots$ boolean;
realit

var
number,
increment
it
: real;
begin. (ind large integral value just beyond integer limits)

ncrement: $: 2$;
While (numbertincrement) = number) do increment: =2*incremont:
adix: :trunc(lnumbertincrement)-number); (



number: $=$ number*radix
end: end of
begin $\left\{\begin{array}{l}\text { fof main program body } \\ \text { find out basic propertios }\end{array}\right.$
onquiry(base, numberofdigits, rounding):

rounding lhen
writelnloutput,' ROUNDED,
Ise writelnloutput,' TRUNCATED')
( compiloute the prection bion bounds)
the
or i:=1 10 numberoldigits do epsilon: =epsilon/base
if
pounding then epsilon:=epsilon/2;


end.
Certification for Burroughs B6700
The following output is produced when running the test program on a Burrough B6748 processor with the University of Tasmania compiler and is correct:

BASE $=$
NUMBER OF
ROMBE OF DIGITS = 13
ROUDED WORST PRECISIONS ARE
$0.9094947 \mathrm{E}-12 \quad 0.7275958 \mathrm{E}-1$

## SOFTWARE TOOLS

One important aspect about Pascal coding style is consistency, although styles certainly differ from one programmer to the next. The two software tools in this issue are both Pascal Prettyprinters, which aid Pascal programmers in their coding activities. They represent 2 vastly contrasting philosophies, and so I think it is appropriate that we
print both, and are assured that we have two of the best in existence. $\mathrm{S}-3$, Prettyprint print both, and are assured that we have two of the best in existence. S-3, Prettyprin adheres to the philosophy that there are serious issues in prettyprinting, and that it is only necessary to impose a minimum set of restrictions in prettyprinting--not be Prettyprint does prettyprinting on a local basis and thus can handle Pascal program fragments, and even incorrect programs. The important principle is that all blank lines and blanks supplied in the original source are preserved.
S. 4 , . Format! indeed does. provide a large set of options because no prettyprinting style can please everyone, and by allowing complete control over the process, one can available, Format is the choice over Prettyprint by 3 to 1 . I use both myself. Prettyprint was first announced in Pascal News 7 etter \#6 page 70, in November, 1976. Henry Ledgard reports that they lost a lot of money distributing it. Charles Fischer we published it.
Format has been around for the last 3 years, and remains in my opinion, one of the alltime, best-looking Pascal programs in existence because of its use of long and meaningful time, best-looking Pascal programs in existence because of its use
identifiers. It looks all the sharper in upper-and-lower case!
There has been quite a bit of noise in the literature about Pascal prettyprinting. We cited Singer, et al,'s article "A Basis for Executing Pascal Programmers" in PUGN 9/10 Sale's article "Stylistics in Languages with Compound Statements" in PUGN 12 page 10 and in this issue; Mohilner's article "Prettyprinting Pascal Programs" in this issue; and I now find Crider's article "Structured Formatting of Pascal Programs" in the November, 1978 SIGPLAN Notices.
Unfortunately, both prettyprinters could do better in their treatment of comments. They are living examples of their results, because they have been run through themselves And as such I am very pleased that we can present them here together with their superb documentation. (*Please excuse my role therein.*) If you want to use these prettyprinters, key them in, or request that your Pascal compiler distributor include them
on the distribution tape for your favorite Pascal system. CDC-6000 Pascal Release 3 will include both Prettyprint and Format. Happy prettyprinting '79!

- Andy Micke1


## S-3 Prettyprint

## Program Title: Pascal Prettyprinting Program

Authors: Jon F. Hueras and Henry F. Ledgard
Computer and Information Science Department
University of Massachusetts, Amherst - August, 1976
Randy Chow and John Gorman).
Bugs corrected by Charles Fischer, Department of Computer Science, University of Wisconsin, Madison.
1977. Indicated by $\lll<$.

Modified for CDC-6000 Pascal Release 3 by Rick L. Marcus University Computer Center, University of Minnesota. 30 September 1978.

## Program Summary:

This program takes as input a Pascal program and
reformats the program according to a standard set of prettyprinting rules. The prettyprinted program is given as output. The prettyprinting rules are given below.

An important feature is the provision for the use of extra spaces and extra blank lines. They may be freely inserted by the user in addition to the spaces and blank lines inserted by the prettyprinter

No attempt is made to detect or correct syntactic errors in the user's program. However, syntactic errors may result in erroneous prettyprinting.

Input File: input - a file of characters, presumably a Pascal program or program fragment.

Output File: output - the prettyprinted program.

## Pascal Prettyprinting Rule

## [ General Prettyprinting Rules ]

1. Any spaces or blank lines beyond those generated by the prettyprinter are left alone. The user is encouraged, for the ake of readability, to make use of this facility
In addition, comments are left where they are found, unless they are shifted right by preceeding text on a line
2. All statements and declarations begin on separate lines.
3. No line may be greater than 72 characters long. Any line longer than this is continued on a separate line.
4. The keywords "BEGIN", "END", "REPEAT", and "RECORD" are forced to stand on lines by themselves (or possibly follwed by upportig comments) "REPEAT-UNTIL" state-
. A blank line is forced before the keywords "PROGRAM" "PROCEDURE", "FUNCTION", "LABEL", "CONST", "TYPE", and "VAR"
5. A space is forced before and after the symbols $":="$ and "=". Additionally, a space is forced after the symbol ":".
Note that only "="s in declarations are formatted. "="s in Note that only "="s in declarati
expressions are ignored.
[ Indentation Rules ]
6. The bodies of "LABEL", "CONST", "TYPE", and "VAR" declarations are indented from their corresponding declaration header keywords.
7. The bodies of "BEGIN-END", "REPEAT-UNTIL", "FOR", "WHILE",
"WITH", and "CASE" statements, as well as "RECORD-END" struc-
"WITH", and "CASE statements, as well as RECORD-END" str
their header keywords.
8. An "IF-THEN-ELSE" statement is indented as follows:

IF <expression>
THEN
ELSE
<statement>

## General Algorithm

The strategy of the prettyprinter is to scan symbols from the input program and map each symbol into a prettyprinting action, independently of the context in which the symbol appears. This is accomplished by a table of prettyprinting
options. options

For each distinguished symbol in the table, there is an associated set of options. If the option has been selected for each option is performed.

The basic actions involved in prettyprinting are the indentation and de-indentation of the margin. Each time the margin is indented, the previous value of the margin is pushed onto a indented: Each time the margin is de-indented, the stack is popped off to obtain the previous value of the margin.

The prettyprinting options are processed in the following order, and invoke the following actions
crsuppress - If a carriage return has been inserted folihowitig the previous symbol, then it is inhibited until the next symbol is printed.
crbefor

- A carriage return is inserted before the aurrent symbol (unless one is already there)
blanklinebefore Aldank line is inserted before the current symbol (unless already there).
dindentonkeys
- If any of the specrfied keys are on top of of the stack, the stack is popped, de-indentimg the margin. The process is repeated suntill the top of the stack is not one of the i) specilied keys.
dindent
- The stack is unconditionally popped and the maxgin -is : de-indented.
spacebefore
- A space iśs insertèa béfore' tiè symbol being scanned (unless already therees).

specialehar $=$ packed array [ $1 . .2$ ] of char;
dblchrset $=$ set of becomes..opencomment;
dblchartable $=$ array [ becomes. opencomment ] of specialchar;
sglchartable $=$ array [ semicolon..period] of char;
string $=$ array [ 1..maxsymbolsize ] of char;
symbol $=\underline{\text { record }}$
ecomes, opencomment, closecomment,
emianong oolon, raq! equails,


option $=($ crsuppress
crbefiorye; in
dindentionikeysi\%,
dindent. :
spacebefiom
spaceafter,
gobbilesymbolis
indentibytath,
indent toolp,
optionset $=$ set of option
keysymset $=$ set orf keysstymbol;
tableentry $=$ record
optionsselected optionset;
, gobblettermindtors: keysymset;
end;
optiontable $=$ array [ keysymboll ] of tableentry;
key $=$ packed array [ 1. maxkeylength ] of char;
keywordtable $=\underline{\text { axray }[\text { progsym. .untilsym ] of key; }}$

```
\begin{tabular}{ll} 
name & : keysymbol \\
valu & : string; \\
length \(\quad\) integer; \\
spacesbefore: & integer;
\end{tabular} spacesbefore: integer;
end;
symbolinfo \(=\) "symbol;
```



```
charinfo \(=\underline{\text { reaord }}\)
valua : charname
end;
        valu : string;
        valu : string;
        cesbefore: integer;
        cesbefore: integer;
        d;
```

        d;
    ```

```

begin { storenextchar }
getchar( { from input }
{updating } nextchar,
{ returning } currchar,)
if length < maxsymbolsize
then
length:= length + 1;
valu [length] := currchar.valu
end
end; { storenextchar }
procedure skipspaces(
{ updating } var currchar,
returrining var nextchar : charinfo;
crsbefore : integer )
begizin { skipspaces }
spacesbefore := 0;
crsbefore :=0;
whtle nextchar.name in_[ blank, endofline] do
Begin
getchar({ from input }
{updating} nextchar,

```
```

if name <> filemark
get(input)
end \{with \}
end; \{ getchar \}
procedure storenextchar ( \{ from input \}
var length
nextchar : charinfo; string );

```

\section*{\(\{\) placing in \} var valu}
```

begin \{ storenextchar

$$
\begin{aligned}
\text { getchar( } & \{\text { from input \} } \\
& \{\text { updating \} nextchar, } \\
& \{\text { returning \} currchar }),
\end{aligned}
$$

then
lengthin: $=$ length +1 ;
valu [length] := currchar.valu
End
procedure skipspaces'
updating \} var currchar,
\{returizing \} Var spacesbefore, $\begin{gathered}\text { crsbefore }\end{gathered}$ : integer );
begizin \{ skipspaces \}
spacesbefore $:=0$;
whitie nextchar.name in [blank, endofline] do
\{updating \} nextchar

```
    case currchar.name of
    case currchar.name of
        blank : spacesbefore := spacesbefore +1 ;
        blank : spacesbefore := spacesbefore +1 ;
        endofline : begin
        endofline : begin
            crsbefore \(:=\) crsbefore +1 ;
            crsbefore \(:=\) crsbefore +1 ;
                    spacesbefore := 0
                    spacesbefore := 0
                end
                end
            end \{case \}
            end \{case \}
end \{while \}
end \{while \}
end; \{ skipspaces \}
end; \{ skipspaces \}
procedure getcomment( \{ from input \}
procedure getcomment( \{ from input \}
                                    \{updating \} var currchar,
                                    \{updating \} var currchar,
                                    \(\begin{aligned} \text { nextchar } & : \text { charinfo; } \\ \text { Var name } & \text { : keysymbol; } \\ \text { Var valu } & : \text { string; }\end{aligned}\)
                                    \(\begin{aligned} \text { nextchar } & : \text { charinfo; } \\ \text { Var name } & \text { : keysymbol; } \\ \text { Var valu } & : \text { string; }\end{aligned}\)
                                    var length \(\quad \begin{aligned} & \text { : integer } \\ & \text { ) }\end{aligned}\)
                                    var length \(\quad \begin{aligned} & \text { : integer } \\ & \text { ) }\end{aligned}\)
begin \{ getcomment \}
begin \{ getcomment \}
name := opencomment;
name := opencomment;
    while not \(\left(\quad\left(\right.\right.\) currchar.valu \(\left.={ }^{\prime} *^{\prime}\right)\) and \(\left(\right.\) nextchar.valu \(\left.\left.\left.=\prime^{\prime}\right) \prime\right)\right)\)
    while not \(\left(\quad\left(\right.\right.\) currchar.valu \(\left.={ }^{\prime} *^{\prime}\right)\) and \(\left(\right.\) nextchar.valu \(\left.\left.\left.=\prime^{\prime}\right) \prime\right)\right)\)
            or (nextchar.name \(=\) endofline )
            or (nextchar.name \(=\) endofline )
        storenextchar ( \{ from input
        storenextchar ( \{ from input
            \{ updating \} length,
            \{ updating \} length,
                                    currchar
                                    currchar
                                    nextchar;
                                    nextchar;
    if (currchar.valu \(=\) '*') and (nextchar.valu \(=\) ' \()^{\prime}\) )
    if (currchar.valu \(=\) '*') and (nextchar.valu \(=\) ' \()^{\prime}\) )
            then
            then
```

hit: boolean;
thiskey: keysymbol;
begin {'idtype }
idtype := othersym;
if length <= maxkeylength
then
for i := 1 to length do
for i := length+1 to maxkeylength do
keyvalu [i] := space;
thiskey := progsym;
hit := false;
while not(hit or (thiskey = succ(untilsym))) do
if keyvalu = keyword [thiskey]
then
hit:= tru
else
thiskey := succ(thiskey);
if hit
idtype := thiskey
end;'
end; { idtype }
.. procedure getidentifier ( { from input }
var currchar,
nextchar : charinfo;
keysymbol;
: string;
integer );
begin {getidentifier }
while nextchar.name in [ letter, digit ] do
storenextchar( { from input }}{$$
\begin{array}{l}{\mathrm{ updating } length,}}
                                currchar,
                                nextchar,
    name := idtype( { of } valu,
    if name 'in [ recordsym, casesym, endsym]
        chase name of 
            recordsym : recordseen := true;
            casesym : if recordseen
```



begin \{ getnextsymbol \}
case nextchar.name of
letter : getidentifier ( \{ from input
nextchar
vame,
);
igit : getnumber \(\{\) from input
nextchar,
name,
length );
quote
: getcharliteral ( \(\left\{\begin{array}{l}\text { from input } \\ \{\text { updating }\}\end{array}
$$\right\}\)
rrchar,
\{returning\} name,
length );
nexttwochars = dblchar [thischar]
then
else
ar := suce(thischar)
if $\frac{\text { not hit }}{\text { then }}$
le not(hit or (pred (thischar) $=$ period)) do
currchar.valu $=$ sglchar [thischar]
hit := true
thischar $:=$ succ(thischar)
hit
$\frac{\text { then }}{\text { chartype }:=}$ thischar
chartype := othersym
```procedure getspecialchar( { from input }             {updating } var currchar,             { returning } \{\mp@code{var name var valu : < <eysymbol;}```
744

gobbleterminators $:=[]$
$\frac{\text { end; }}{\frac{\text { with ppoption }[\text { typesym }]}{\text { begin }} \text { optionsselected }}:=$
blanklinebefore, dindentonkeys, spaceafter,
dindentsymbols $:=[$ labelsym,
gobbleterminators $:=[]$
with ppoption [ varsym] do
begin
optionsselected := [ blanklinebefore, dindentonkeys, spaceaf ter,
dindentsymbols $:=[$ labelsym
typesym ];
gobbleterminators := []
end;
with ppoption [ beginsym] do
optionsselected $:=$ [ dindentonkeys crafter ];
dindentsymbols $:=[$ labelsym, constsym, typesym,
gobbleterminators := []
with ppoption [ repeatsym] do begin
optionsselected $:=[$ indentbytab
號 $:=[]$
end;
with ppoption [recordsym] do $\frac{\text { begin }}{\text { op }}$
optionsselected $:=$ [ indentbytab,
dindentsymbols $:=[] ;$
end;
wi.th ppoption [ casesym] do begtin
optionsselected := [ spaceafter, ndentbytab, crafter ];

$$
\text { dindentsymbols } \quad:=[7 ;
$$

$$
\text { gobbleterminators }:=[\text { ofsym }]
$$

endes
witith ppoption [ casevarsym ] do begin

| -1009 | opEionsselected := [ spaceafter, |
| :---: | :---: |
| 1010 | indentbytab, |
| 1011 | gobblesymbols, |
| 1012 | crafter ]; |
| 1013 | dindentsymbols := []; |
| 1014 1015 | gobbleterminators $:=$ [ ofsym ] |
| 1016 |  |
| 1017 | with ppoption [ ofsym ] do |
| 1018 | begin |
| 1019 | optionsselected := [ crsuppress, |
| 1020 | spacebefore ]; |
| 1021 | dindentsymbols gobbleterminators $:=[1]$ |
| 1023 | end; |
| 1024 |  |
| 1025 | with ppoption [ forsym ] do |
| 1026 | begin |
| 1027 | optionsselected := [ spaceafter, |
| 1028 1029 | indentbytab, |
| 1030 | crafter ]; |
| 1031 | dindentsymbols := []; |
| 1032 | gobbleterminators := [ dosym ] |
| 1033 | end; |
| 1034 |  |
| 1035 | with ppoption [ whilesym ] do |
| 1036 | begin |
| 1037 1038 | optionsselected := [ spaceafter, |
| 1038 1039 | indentbytab, gobblesymbols |
| 1040 | crafter ]; |
| 1041 | dindentsymbols $:=$ []; |
| 1042 | gobbleterminators := [ dosym] |
| $\begin{aligned} & 1043 \\ & 1044 \end{aligned}$ | end; |
| 1045 | with ppoption [ withsym] do |
| 1046 | begin |
| 1047 | optionsselected $:=$ [ spaceafter, |
| 1048 | indentbytab, |
| $\begin{aligned} & 1049 \\ & 1050 \end{aligned}$ | gobblesymbols, |
| 1051 | dindentsymbols $:=[] ;$; |
| 1052 | gobbleterminators $:=[$ dosym ] |
| 1053 | end; |
| 1054 |  |
| 1055 | with ppoption [ dosym ] do |
| 1056 | begin |
| 1057 | optionsselected := [ crsuppress, |
| 1058 | [spacebefore ]; |
| 1059 | dindentsymbols := []; |
| 1060 | gobbileterminators := [] |
| 1061 1062 | end; |
| 1063 | with ppoption [ ifsym] do |
| 1064 | begin |
| 1065 | optionsselected := [ spaceafter, |
| 1066 | indentbytab, |
| 1067 | gobblesymbols, |
| 10 |  |
| 4069 | dinidentsymbols $:=$ []; |
| 1070 1071 | gobbleterminators $:=$ [ thensym] |
| 1072 | , |
| 1073 | Trith ppoption [ thensym] do |
| 1074 | begin |

ith ppoption [ becomes ] do. begin
[ spacebefore, spaceafter,
$\begin{aligned} \text { dindentryymols }:= & {[1 ;} \\ \text { gobbleterminators }:= & {\left[\begin{array}{l}\text { endsyn, } \\ \\ \text { untilsym, }\end{array}\right.}\end{aligned}$


With ppoption [ untilsym] do

| $\frac{\text { with }}{\text { begin }}$ ppoption [untilsym] do |  |
| :---: | :---: |
| optionsselecteddindentsymbols | := [ crbefore, |
|  | dindentonkeys, dindent, |
|  | spaceafter, |
|  | gobblesymbols, |
|  | $\begin{aligned} & \text { crafter ]; } \\ & \text { ifsym, } \end{aligned}$ |
|  | thensym, |
|  | elsesym, |
|  | fors ym, |
|  | whilesym, |
|  | witithsym, |
| gobbleterminators | colon, equals ]; |
|  | := [ endsym, |
|  | untilsym, |
|  | elsesym, semicolon |
| end; |  |
| with ppoption [ becomes ] do , begin |  |
|  |  |
| optionsselected | := [ spacebefore, spaceafter, gobblesymbols ] |
| . di.xdentsymbols gobbleterminators | $t=1]$; |
|  | := [ endsym, |
|  | untilsym, |



```
end; \{insertcr \}
340
341
1342 procedure insertblankline( \{ updating \} \(\begin{aligned} & \text { var } \\ & 1343 \\ & \{\text { writing to output \} }\end{aligned}\)
const
        once \(=1\);
1347
1348
1348
1349
begin \{ insertblankline \}
if currsymn \({ }^{\wedge}\).crsbefore \(=0\)
        \(\xrightarrow[\text { beg }]{ }\)
            if currlinepos \(=0\)
                then
                    \(\frac{\text { writecrs ( once, },\{\text { updating \} currlinepos }}{\{\text { writing to output \} }}\)
                else
                    -writecrs ( twice, \(\left\{\begin{array}{l}\text { updating \} currlinepos } \\ \{\text { writing to output }\}\end{array}\right.\) );
            currsym².spacesbefore :=0
            end
        else
        if currsym².crsbefore \(=1\)
            \(\frac{\text { then }}{\text { if }}\) currlinepos \(>0\)
                    then
                            writecrs (once, \{updating \} currlinepos ,
end; \{ insertblankline \}
procedure lshifton( \{ using \} dindentsymbols : keysymset );
var
    indentsymbol: keysymbol
        prevmargin : integer;
begin \{ lshifton \}
if not stackempty
        \(\frac{\text { then }}{\text { begi }}\)
            repeat
                    popstack( \{ returning \} indentsymbol,
                    if indentsymbol in dindentsymbols
                    then
                currmargin := prevmargin
                    until not(indentsymbol in dindentsymbols)
                    오 (stackempty);
                    if not(indentsymbol in dindentsymbols)
```

begin \{insertspace \}
if currlinepos < maxiinesize
then

## write(output, space);

currlinepos $:=$ currlinepos +1 ;
$\left.\frac{\text { with symbol }}{\text { if }(\text { do }}=0\right)$ and $($ spacesbefore $>0)$

## end

## end; \{insertspace \}

```

procedure movelinepos( \{ to \} newlinepos : integer;
        \(\left\{\begin{array}{l}\text { newlinepos } \\ \text { from } \\ \text { far } \\ \text { currlinepos }\end{array}\right.\)
        writing to output
\(\underline{\mathrm{Var}}_{i ;}\) integer;
begin \{ movelinepos \}
for \(\mathrm{i}:=\) currlinepostl to newlinepos do
                    write(output, space);
1467 for \(i \quad:=\) currlinepos +1
1468 to newlinepos do
1470 currlinepos := newlinepos
```

1465

- 5

```
end: {movelinepos }
```




```
var i: integer;
begin { printsymbol }
    with currsym^ do,
            for i := 1 to lergth do
                    write(output, valu[i]);
                startpos := currlinepos; { save start pos for tab purposes } <<<<,
            currlinepos := currlinepos + length
        end {with}
end; { printsymbol }
procedure ppsymbol({ { in } currsym - : symbolinfo
Const once = 1
var}n\mathrm{ newlinepos: integer
begin { ppsymbol }
    with currsym^ do
        begin
            writecrs( {using } crsbefore,
                    {updating } currlinepos
            if (currlinepos + spacesbefore > currmargin)
                    or (name in [ opencomment, closeconment ])
                    then
                    else
                    newlinepos := currnargin
            if newlinepos + length > maxlinesize
                    then
                    writecrs( once, { updating } currlinepos
                            );
                    if currnargin + length}<=\mathrm{ maxlinesize
                    then}\mathrm{ newlinepos := curmnargin
                    else
                    ewlinepos := currnargin
```



```
        end;
    movelinepos( { to } newlinepos,
    from} currlinepos
    {in output }
    printsymbol( { zn } currsym,
        ingring.cecurrsym,
    end {with }
end; { ppsymbol }
procedure rshifttoclp( { using } currsym : keysymbol );
    forward;
procedure gobble( { symbols from input }
            (up to
            updating } var currsy
            {writing to output } nextsym : symbolinfo
begin { gobble }
    rshifttoclp( { using } currsym^.name );
    while not(nextsym^.name in (terninators + [endoffile])) do
        begin
            getsymbol( ( { from input } 
            ppsymbol({ { in } currsym ,
        end; { while }
    lshift
end; { gobble }
procedure rshift( { using } currsyn : keysymbol );
begin {rshift }
    if not stackfull
            then
                currmargin);
    { if extra indentation was used, update margin. }
    if startpos > currmargin
    if startpos > currmargin
```



[^0]
if indentbytab in optionsselected
$\frac{\text { chen }}{\text { rshifte }}$ ( using \} currsymn.name );
if indenttoclp in optionsselected
$\xrightarrow{\text { then }}$
if gobblesymbols in optionsselected gob
obble( \{ symbols from input \}

$\left\{\begin{array}{l}\text { up to }\} \\ \text { updating \} } \\ \text { currsym, }\end{array}\right.$ nextsym
\{writing to output \} );
if crafter in optionsselected
$\frac{\text { then }}{\text { crpending }}:=$ true
end \{with \}
end; \{while \}
if crpending
$\frac{\text { then }}{\text { writeln(output) }}$
end.
pascal program formatter
S-4 Format

$$
\begin{aligned}
& \text { - Michaej N. Condict } \\
& \text { - Rick. L. Marcus } \\
& \text { - Andy Mickel }
\end{aligned}
$$

What Format Does

Format is a flexible prettyprinter for Rascal programs. It takes as input a syntactically-correct Raseal program and prodtces as oatput an equivalent but reformatted
Pascal programe: The nesulting progqam eonsists of the same sequence: and comments, but they, are rearnanged, with respect to line boundaries and columins for
readability. readabili.ty -
structurfes matintains monsistent spaging between newmbols, breaks control and data
 line-numbers añ automatic oomments, or deleting madi unpegessary blanks to saye space, are described bedow.

The flexibility of Format is accomplished by allowing you to supply various directives (options) which override the default values. Rather than being a rigid prettyprinter which decides for you how your program is to be formatted, you have the ability to control how formatting is done, not only prior to execution but also during
execution through the use of prettyprinter directives embedded in your program.

Experience with Format over the last three years has shown that most users can find
set of values for the directives which produce satisfactory results. The defaul values are typical.

How To Use Format

The use of Format will vary from implementation to implementation, but will involve one major input file containing a Pascal program and one output file for the reformatted program. Additionally it may be possible to supply the initial values of directives to Format when it begins execution.

Directives to Format may always be specified in the program itself inside comments with a special syntax. Thus the first line of a program is an ideal spot for a comment containing directives. Subsequent use of embedded directives allows you to change the comments is given below (The syntax is given using "EBNF"--Extended Backus-Naur Form--see Communications ACM, November, 1977, page 822.):

DirectiveGomment = "(*" DirectiveList "*)"|
"(*\$" CompilerOptionList CommentText DirectiveList "*)".
DirectiveList = "[" Directive \{"," Directive\} "]" CommentText.
Directive $=$ Letter Setting.

Setting $=$ Switch $\mid$ Value $\mid$ Range.
Switch = "+" | "_" .
Value = "=" UnsignedInteger.
Range = "=" UnsignedInteger "-" UnsignedInteger ["<" | ">"].
UnsignedInteger $=$ Digit\{Digit $\}$.
CommentText $=$ \{Any character except " $\}$ " or close-comment $\}$.
Note: As defined above, a Directive may be within a comment specifying a Pascal CompilerOptionlist. On most implementations this is a "\$" followed by a series of Examples of DirectiveGomments:
$(*[A=15, E=3, N=1,1<] *)$ - good for publication quality.
$(*[G=0, W=1-100, C+] *)$ - good for compact storage.
(*\$U+ $[R=1-72, I=2] *$ ) - an example of a Directivelist with a CompilerOptionList.

## Directives to Format

## A=n Align declarations

The A directive forces the alignment of ":" and "=" in declarations. If $A$ is
set to a value greater than 0 , then $n$ should be equal to the maximum identifier
length for that section of your program. The A directive visually clarifies the declaration part of your program. See example below.
Default: $\mathrm{A}=0$ (no alignment).
$B+$ or $B-\quad$ Bunch statements and declarations reasonably.
B+ will place as many statements or declarations onto one line as will fit $\triangle$ within the specified write margins (W directive) subject to readability $\Gamma$ within the specified write margins constraints. Bunching (B+) when the display is off ( $\mathrm{D}-$ ) has no effect. In $z$ general, B+ saves paper and prevents your program from becoming overly stretched in the vertical direction. See example below.
Default: B- (one statement or statement part per line)
C+ or C- Fully Compress program.
C+ removes all non-essential blanks, end-of-1ines, and comments from your program. A compilable, packed program will be written within the specified $W$ program. A comp
write margins (W directive). The number of spaces specified by the $G$ directive will still be written between symbols. C+ might save some space on long-term storage media such as disk; you might store a program in compressed form and expand it later by reformatting with $C$ -
Default: C-.
D+ or D- Turn Display on or off
D allows you to selectively display portions of your program during formatting. Therefore, D must be switched on and off with directives which are appropriately publication (such as one or more procedures) without having to print the whole program.
Default: D+.
E=n Supply END comments.
The E directive generates comments after "END" symbols if none are already $\theta$ there. Common Pascal coding styles frequently employ these comments. E=1 m
creates comments after the "END" symbol in compound statements which are within structured statements, as well as those constituting procedure and function bodies. The comments take the form: (*StatementPart*) or (*ProcedureName*). 3 $\mathrm{E}=2$ creates comments after the "BEGIN" and "END" symbols constituting procedure 3 and function bodies only. $\mathrm{E}=0$ creates no comments at all. $\mathrm{E}=3$ means $\mathrm{E}=1$ and $\quad$ $\mathrm{E}=2$. See example below.
Default: $\mathrm{E}=2$.
F+ or F- Turn Formatting on or off.
F allows you to format selected portions of your program. F- causes Format to $\omega$ copy the input program directly with no changes. Therefore by switching $F$ on and off with directives which are appropriately placed in your program, you can $C_{\infty}$ preserve text which is already properly formatted (such as comments).
Default: F+ (of course!).
$G=n \quad$ Specify symbol Gap.
The $G$ directive determines the number of spaces placed between Pascal symbols during formatting. $G=0$ still places one space between two identifiers and reserved words. The symbols [ ] (), and : are handled independently of $G$.
Default: $G=1$.
I=n Specify Indent tab
I indents each nesting level of statements and declarations a given number of columns. Using $\mathrm{I}=2$ or $\mathrm{I}=1$ helps prevent excessively-narrow lines within the specified write margins (W directive) where there are heavily-nested constructs. Default: $I=3$.
$\mathrm{L}=\mathrm{n}$ Specify Line-wraparound indent tab.
determines the indentation of the remainder of statements or declarations
L determines the indentation of the
which are too long to fit on one line.
Default: $L=3$.
$N=x-y<$ or $N=x-y>$ Generate line-numbers on the left or right.
The $\mathbb{N}$ directive indicates the starting line-number ( $x$ ) and the increment ( $y$ ) for
each succeeding line－number．If $\mathrm{y}>0$ then line－numbers are written outside the specified write margins for the formatted program in sequential order starting at $x ; y:=0$ shuts off line－numbering．＂く＂writes up to 4 －digit，right－justified line－numbers together with a trailing space to the left of each linie．＂＞＂ writes 6－digit，zero－filled line－numbers to the right of each line．Use the $\mathbb{N}$ directive along with the $W$ directive
$\mathrm{P}=\mathrm{n}$ Specify spacing between Procedure and function declarations．
The $P$ directive determines the number of blank lines to be placed between procedure and function declarations．$n>2$ makes procedures and functions visually stand out
Default：$P=2$
$\mathrm{R}=\mathrm{x}-\mathrm{y}$ Specify Read margins．
The R directive indicates which columns are significant when Format reads from The R directive indicates which columns are significant when Format reads from
input fille．$R$ allows Format to accept files which have line－numbers in the finput（ $x-1$ ）columns or after the yth column．
Default： $\mathrm{R}=1-999$（large enough to read to end－of－line in most cases）．
$S=n \quad$ Specify Statement separation．
The $S$ directive determines the number of spaces between statements bunched on the same line by the use of the B＋directive．Note that this directive is only in elfect if B＋is used．
Default：$S=3$ ．
W＝x－y Specify Write margins
The $W$ directive indicates which columns are used for writing the reformatted program on the output file．Any line－numbers generated（N directive）are written outside these margins．
Default：$W=1-72$ ．

## Examples

The A Directive

Here is a sample program fragment before using Format
PROGRAM SAMPLE（OUTPUT）；
CONST $A=6$ ；： $\mathrm{AB}={ }^{\prime}=$ LETETERS ${ }^{\prime}$ ；THREE $=3$ ；
TYPE RANGE＝1．．6；
$\mathrm{COLOUR}=(\mathrm{RED}, \mathrm{BLOE})$ ；
I，I2，I33，I444，I5555：RANGE；
YES，NO；MAYBE：BOOLEAN；
BEGIN END．
Here is the output from Format with all defaults set：
PROGRAM SAMPLE（OUTPUT）；
${ }^{\text {CONSI }}=$
$A=6$ ；


- － $\operatorname{THREE}=3$ ；

TYPE ：er ：：
RANGE $=$
COLOUR $=$
（RED，BTUEX ；＂


YES，NO，MAYBE：BOOLEAN；
begin
END（＊SAMPLE＊）．
Here is the output from Format having added a line with the $A=5$ directive
（＊［A＝5］ALIGN DECLARATIONS．＊）
PROGRAM SAMPLE（OUTPUT）；
CONST

$$
A=6 ;
$$

$$
\begin{aligned}
A & =6 ; \\
A B C & =\text { 'LETTERS' }
\end{aligned}
$$

$\begin{aligned} A B C & =1 \\ \text { THREE } & =3 ;\end{aligned}$
TYPE
$\begin{aligned} \text { RANGE } & =1 \ldots 6 ; \\ \text { COLOUR } & =(\text { RED，BLUE }) ;\end{aligned}$
var
$I$,
$I 2$,
I33，
I444，
I444， I5555：RANGE；
YES，
MAYBE：BOOLEAN；
BEGIN
END（＊SAMPLE＊）．，
The B Directive
If the input to Format is：
PROGRAM T（OUTPUT）；
CONST INCREMENT $=5$ ；
VAR I，J，N：INTEGER；
BEGIN
$\mathrm{N}:=0$ ；
$\mathrm{J}:=3 ; \mathrm{I}:=\mathrm{SQR}(\mathrm{N}) ; \mathrm{N}:=\mathrm{N}+\mathrm{INCREMENT} ;$
IF N $>73$ THEN BEGIN DOTHIS；DOTHAT END
IF N＞5 THEN IF J＞6 THEN $\quad$ DOSOMETHINGELSE；
END．
then the output from Format（using the default，B－）is：
PROGRAM T（OUTPUT）；
CONST
INCREMENT $=5$ ；
${ }^{\text {VAR }}$ I，J，N：INTEGER；
BEGIN

$$
\begin{aligned}
& \begin{array}{l}
N:=0 \\
\mathrm{~J}:=3 \\
\mathrm{I}:=\mathrm{S} \mathrm{Q} R(\mathrm{~N})
\end{array} \\
& \mathrm{N}:=\mathrm{N}+\text { TNCREMENT; }
\end{aligned}
$$

$$
\begin{aligned}
& \text { Dotitis; }
\end{aligned}
$$

```
", IF \(\operatorname{END}>5\) THEN
    IF J > \(>6\) THEN
TDOSOMEHHINGESE
```


and the :output from Format with B directives embedded is:
(*[BHIT:BUNOH:STATEMENTS. *)
PROGRAM.T(OUGEUR);
cons'T
INCREMENT $=5$;
VAR
I; Js.N: INTEGER:
BEGIN
$N:=0 \quad J:=3 k \quad I \quad=\operatorname{SQR}(N) ; \quad N:=N+$ INCREMENT;
( IF F N S .73 . THEN BEGTN DOTHIS; DOTHAT END;
(*[B-] UNBUNCH. *)

IF DOSOMETHENGELSE;
END (*T*).
The E Directive

Suppose that a Pascal program fragment looked like

$$
\begin{aligned}
& \text { PROCEDURE SAMPLE; } \\
& \text { PROCEDURE INNER; } \\
& \text { BEGIN END; } \\
& \text { BEGIN } \\
& \text { IF } X=3 \text { THEN: } \\
& \text { BEGIN } X:=1 ; I:=I+1 \\
& \text { END } \\
& \text { ELSE. } \\
& \text { BEGIN } X:=X+I ; I:=0 \\
& \text { END; } \\
& \text { WHILE (CH<> } X^{\prime} \text { ) AND FLAGI DO } \\
& \text { BEGIN } I:=I+3 \text {; INNER END; END; }
\end{aligned}
$$

then using Format with E=3 produces:

## PROCEDURE SAMPLE;

```
PROGEDURE INNER;
    BEGIN
    END (*INNER*);
```

BEGIN (*SAMPLE*)
IF $X=3$
$\underset{\text { THEN }}{\text { IF }} \mathbf{X}=3$
THEN BEGIN
$X:=1 ;$
$I:=I+$
$\operatorname{END}^{I}:=I+1$
ELSE
ELSE BEGIN
$X:=X+I$
$I:=0$


## How Format Works

Format parses your program by performing syntax analysis similar to the Pascal compiler: fectussive descent within nested declarations and statements. It gathers characters into a buffer in which the indenting count of each character is maintained. The characters are being continually emptied from the buffer as new ones are added.

Format has limfted error-recovery facilities, and no results are guaranteed if a syntactically-incorrect program is supplied as input.

The bane of most Pascal prettyprinters is the treatment of comments. Format considers them in the context of a decianation or statement. Therefore using comments like:

CONST LS $=6$ (*LINESIZE*);
is a good idea because format will carry the comment along with the declaration. Similarly:

$$
\text { BEGIN (* ' } Z^{\prime}<\mathrm{CH}<={ }^{\prime} \text { '*) }
$$

is also okay.
Stand-alone comments however, receive rough treatment from Format. The first line of such comments are always left-justified and placed on a separate line. See the $F$ irective. Thus:

CONST LS=6; (*LINESIZE*)
will be reformatted as:

> CONST LS = $6 ;$ (*INESIZE*)

Proper treatment of comments is certainly an area of future development for Format.
Format issues the following error messages:

1. " *** 'PROGRAM' EXPECTED.'

The Pascal program you fed to Format did not contain a Standard Pascal program declaration.
2. " *** ERRORS FOUND IN PASGAL PROGRAM."

Your program is syntactically incorrect. The output from Format probably does not contain all of the text from your input file. The cause could be any syntactic error, most commonly unmatched "BEGIN-END" pairs, or the lack of semicolons, string quotation marks, or the final period.
3. " *** STRING TOO LONG."

Your program contains a character string (including both the quotes) which is wider than the specified write margins (W directive).
4. " *** NO PROGRAM FOUND TO FORMAT."

The input file given to Format is empty.

## History

Format was originally written in 1975 by Michael Condict as a class project in a Pascal programming course taught by Richard Cichelli at Lehigh University using CDC-600 Pascal. After that, making improvements and adding directives became, temporarily, an
obsession with the author (note limited usefulness of the D directive). Fortunately, the program eventually stabilized and is now in general use by Pascal programmers at Lehigh University and other institutions. After graduation the author transported Format in 1977 to a PDP-11 running under the Swedish Pascal compiler and RSX-11 with a total effort of 2 days.

Assistance in bringing up Format may be obtained by writing to Michael Condict at Pattern Analysis and Recognition Corporation, 228 Liberty Plaza, Rome, NY 13440. Forma has been made as portable as possible, but portable programs are hampered by non-standard
character sets and non-standard techniques for associating program objects (e.g. file variables) with operating system objects (e.g. files as mass-storage devices).

The PDP-11 version of Format uses a procedure
which serves a function similar to standard Pascal program headers for external files This version accepts initial values for directives after it types a prompt for you at your interactive terminal.

On the other hand, the CDC-6000 version accepts initial values for directives after a "/" on the operating system control statement which executes Format.

Format was modified for inclusion with Release 3 of CDC-6000 Pascal by Rick L. Marcus and Andy Mickel, University Computer Center, University of Minnesota, in November, 1978

\section*{\{ $\mathrm{A}=20, \mathrm{~B}+, \mathrm{R}=1-100, \mathrm{I}=2, \mathrm{~S}=2]$ FORMATTER DIRECTIVES. <br> \{ <br> | * | PASCAL PROGRAM FORMATTER |  |
| :---: | :---: | :---: |
| * |  |  |
| * | AUTHOR: MICHAEL N. CONDICT, 1975. |  |
| * |  |  |
| * |  | LEHIGH UNIVERSITY |
| * | CURRENT ADDRESS: | Par corp. |
| * |  | 228 LIBERTY PLė̇ZA |
| * |  | ROME, NY 13440 |
| * |  |  |
| * | UPDATED: AUGUST, 1978. |  |

${ }^{3}$
program Format(Input, Output);
$\frac{\text { Iabel }}{13 ;}$
const
AlfaLeng $=10$;
\{ !!!!!!! IMPLEMENTATION DEPENDENCY: !!!!!!!!
MAXIMUM Char VALUES. $\}$
MinChar $=0 ;$
MaxChar $=127$

$$
\begin{aligned}
\text { MaxChar } & =127 ; \\
\text { LastPascSymbol } & =29 ;
\end{aligned}
$$

\{ the following Constants Must all be changed together, so that thein VALUES AGREE WITH THEIR NAMES:
2 \} BufferSize $=160 ;$

$$
\left.\begin{array}{rl}
\text { BuffSzPI } & =161 ; \\
\text { BuffSzM1 } & =159 ; \\
\text { BuffSzDiv10 } & =16 ; \\
\text { MaxReadRightCol } & =999 ; \\
\text { MaxWriteRightCol } & =72 ;
\end{array}\right\}
$$

type
Alfa $=$ packed array [1 .. AlfaLeng] of Char ( SET STI IMPLEMENTATION DEPENDENGY: !!!!!!!!\}

MAY NOT ALLOW SET OF Char.
CharSet = set of Char;
StatmntTypes $=$ (ForWithWhileStatement, RepeatStatement,
IfStatement, CaseStatement, CompoundStatement,
Symbols $=$ (ProgSymbol, Comment, BeginSymbol, EndSymbol,
Semicolon, ConstSymbol, TypeSymbol,
RecordSymbol, ColonSymbol, EqualSymbo1,
PeriodSymbol, Range, CaseSymbol, OtherSymbol,
IfSymbol, ThenSymbol, ElseSymbol, DoSymbol,
OfSymbol, ForSymbol, WithSymbol, WhileSymbol,
RepeatSymbol, UntilSymbol, Identifier,
VarSymbol, ProcSymbol, FuncSymbol, LeftBracket,
RightBracket, CommaSymbol, LabelSymbol,
Width $=0 \ldots$ Buffersize;
Margins $=-100 \ldots$ BufferSize;
SymbolSet = set of Symbols;
OptionSize $=-99.99$
CommentText $=$ array $[1 .$. BuffSzDivi0] of Alfa;
SymbolString = array [Width] of Char;

## ChIsEOL, NextChIsEOL: Boolean;

I: Integer \{USED AS FOR LOOP INDEX\};
ReadCacter: Char;
ReadRightCol: $0 . .1000$;
OutputCol,
WriteColumn,
Actuallef tMargin,
ReadLeftCol,
ReadLef tCol,
WriteLeftCol,
WriteRightCol: Margins;
DisplayIson,
ProcNamesWanted,
ndCommentsWanted,
PackerIsOff,
SavedBunch,
BunchWanted,
NoFormatting: Boolean;
LineNumber,
Increment: Integer;
IndentIndex,
LongLineIndent,
DeclarAlignment,
StatmtSeparationt
ProcSeparation: OptionSize
LastSymbol,
SymbolName: Symbols;
AlphaSymbois,
EndLabel,
EndConst,

```
            EndType,
            EndVar: SymbolSet;
            Symbol: SymbolString;
            Length: Width;
        Length:
LastProgPartWasBody: Boolean;
            Digits,
    LettersAndDigits: CharSet;
            Oldest: Width;
                arCount: Integer { COUNT OF TOTAL CHARS READ,
                    BUT IS OFF BY BufferSize AFTER END OF FIRST BODY
                    IT IS IMPERATIVE THAT CharGount BE MONOTONICALLY
                    INCREASING DURING PROCESSING OF A BODY, AND THAT IT
                    NEVER RETURN TO A VALUE <= BufferSIze, AFTER PASSING
                THAT POINT. THUS "DoBlock" MAY RESET IT AS LOW AS
                "MoxTt - buffrS THL" CHPACTERS }
            e" CHARACTERS. };
        MainNmLength: Width;
            Blanks,
            Zeroes: Alfa;
LettersAndDytits: Charset
CharCount: Width
BUT IS OFF BY Buffersize AFTER END OF FIRST BODY. INCREASING DURING PROCESSING OF A BODY, AND THAT IT
NEVER RETURN TO A VALUE <= BufferSize, AFTER PASSING THAT POINT. THUS "DoBlock" MAY RESET IT AS LOW "MaxInt - Buffersize" Characters. \};
MainNmLength: Width;
Blanks,
Zeroes: Alf
```

            UnWritten: array [Width] of record
                    Ch: Char;
                    ChIsEndLine: Boolean
                        IndentAfterEOL: Margins
        PascalSymbol: array [1 .. LastrascSymbol] of Alfa;
        ascSymbolName: array [1... LastPascSymbol] Of Symbols;
    StatementType0f: array [Symbols] of StatmntTypes;
    procedure ConstantsInitialization
$\frac{\text { begin }}{\text { Main }[1]}:=$ 'MAIN $\quad$; MainNmLength $:=4$
Main [1] $:=$ 'MAIN $\quad$ '; MainNmLength $:=4 ;$
Blanks $:=$;
for $I:=0$ to Buffszml do
with UnWritten[I] do
begin Ch $:={ }^{\prime} \mathrm{A}^{\prime} ;$ ChIsEndLine $:=$ False; IndentAfterEOL $:=0$
end;

NameOf [Character] := OtherSymbol;
bol; (MaxChar) do
Character : $=$, '; NameOf [' ('] $:=$ LeftParenth
NameOf [')']:=RightParenth; NameOf ['='] := EqualSymbol;
NameOf $\left[\prime^{\prime}, \prime^{\prime}\right]:=$ CommaSymbol; NameOf $\left[\prime^{\prime}{ }^{\prime}\right]:=$ PeriodSymbol
NameOf ['['] := LeftBracket; NameOf[']'] := RightBracket;
NameOf $\left['^{\prime}: \prime^{\prime}\right]:=$ ColonSymbol; NameOf $\left[\prime^{\prime}<^{\prime}\right]:=$ EqualSymbol;
NameOf $\left[\prime^{\prime}>^{\prime}\right]:=$ EqualSymbol; NameOf $\left[{ }^{\prime} ;{ }^{\prime}\right]:=$ Semicolon;
PascalSymbol[1] $:=$ 'PROGRAM $;$ PascalSymbol[2] $:=$ 'BEGIN
PascalSymbol[3]:='END
PascalSymbol[5] := 'TYPE
PascalSymbol[7]:='RECORD
PascalSymbol $[9]:=$ 'IF
PascalSymbol $[11]:=$ 'ELSE
PascalSymbol[13] := 'OF
PascalSymbol[15]:= 'WHILE
PascallSymbol 117$]:=$ ' REPEAT
Pasalsymbol 17$]:=$ REPEAT
PascalSymbol $[19]:=-$ PROCEDURE
Pascalsymbol[19] $:=$ 'PROCED
PascalSymbol 21$]:==$ IABEL
PascalSymbol[23] $:=$ 'LABE
PascalSymbol[25]:='AND
PascalSymbol [27] $:=$ 'NOT
PascalSymbol[29] $:=$ 'NOSYMBOL
$\begin{array}{lll}\text { PascalSymbol }[27] & :=\text { 'NOT } & ; \\ \text { PascalSymbol[29] } & :=\text { PascalSymbol[28] } & :=\text { 'ARRAY } \\ \text { PascSymbolName }[2] & :=\text { BeginSymbol; } ; & \text { PascSymbolName[1] }:=\text { ProgSymbol; } \\ \text { PascSymbolName [3] } & :=\text { EndSymbol; }\end{array}$
$\begin{array}{ll}\text { PascSymbolName [2] }:=\text { BeginSymbo1; } & \text { PascSymbo1Name [3] }:=\text { EndSymbol; } \\ \text { PascSymbolName [4] } & :=\text { ConstSymbol; }\end{array}$
$\begin{aligned} & \text { PascalSymbol }[2] \\ & \text { PascalSymbol }[4]\end{aligned}:=$ 'BEGIN
PascalSymbol[4]:='CONS
PascalSymbol[6]:='VAR
PascalSymbol[8]:= 'CASE
PascalSymbol[10] $:=$ 'THEN
PascalSymboi[12] $:=$ 'DO
PascalSymbol[12] $:=$ 'DO
PascalSymbol[14] $:=$ 'FOR
PascalSymbol[14] $:=$ 'FOR
PascalSymbol[16] $:=$ 'WITH
PascalSymbol[18] $:=$ 'UNTIL
PascalSymbol[20] $:=$ 'FUNCTION
PascalSymbol[20]:='FUNCTION
PascalSymbol[22] $:=$ 'IN
PascalSymbol[24]:='DIV
PascalSymbol[26] $:=$ 'OR
PascalSymbol[28] $:=$ 'ARRA

UnWritten: array [Width] of record
Ch: Char;
ChIsEndLine: Bool IndentAfterEOL: Margins

Name0t: array [Char] of Symbols;
StatementType0f: array [Symbols] of StatmntTypes;
procedure ConstantsInitialization;
$\frac{\text { begin }}{\text { Mai }}$
Main $[1]:=$ 'MAIN $\quad$ '; MainNmLength $:=4 ;$
Blanks $:=$,
for $I:=0$ to Buffszml do
begin $\mathrm{Ch}:={ }^{\text {th }} \mathrm{A}^{\prime} ;$ ChIsEndLine $:=$ False; IndentAfterEOL $:=0$ end;
NameOf [Character] := OtherSymbol;


$\begin{array}{ll}\text { PascalSymbol[1] }:=\text { 'PROGRAM } & \text {; PascalSymbol[2] } \\ \text { PascalSymbol }[3]:=\text { 'BEGIN } \\ \text { 'END } & \text {; PascalSymbol }[4] \\ :=\text { 'CNST }\end{array}$
PascalSymbol[5]:='TYPE
PascalSymbol[7] $:=$ 'RECORD
PascalSymbol[11] $:={ }^{\prime}$ ELSE
PascalSymbol[13] $:={ }^{\prime}$ OF
PascalSymbol[15] $:=$ 'WHILE
PascalSymbol 177$]:=$ REPEAT
PascalSymbol[19] :='PROCEDURE'
PascalSymbol[21] $:=$ 'LABEL
PascalSymbol $[23]:=$ 'MOD
PascalSymbol[25]:='AND
$\begin{array}{ll}\text { PascSymbolName [2] }:=\text { BeginSymbol; } & \text { PascSymbo1Name [3] }:=\text { EndSymbol; } \\ \text { PascSymbolName [4] } & :=\text { ConstSymbol; }\end{array}$

PascSymbolName[6] := VarSymbol; PascSymbolName [7] := RecordSymbol;
PascSymbolName [8] := CaseSymbol; PascSymbolName [9] := IfSymbol;
PascSymbolName [10] := ThenSymbol; PascSymbolName [11] := ElseSymbol;
PascyymbolName[12] $:=$ DoSymbol; PascSymbolName [13] $:=$ OfSymbol;
PascSymboIName [16] := WithSymbol;
PascSymbolName[17] := RepeatSymbol;
PascSymbolName[18]:=UntilSymbol;
PascSymbolName [21] := Labe1Symbol;
PascSymbolName [29]:= Identifier;
for $I$ := 22 to 28 do PascSymbolName [I] := AlphaOperator;
for SymbolName $:=$ ProgSymbol to AlphaOperator do
StatementTypeof [SymbolName] $:=$ otherStatement;
StatementTypeof [BeginSymbol] $:=$ CompoundStateme
StatementTypeof [CaseSymbol] $:=$ CaseStatement;
Statementypeof [Ifsymbol] := IfStatement;
StatementTypeOf [ForSymboil]:= ForWithWhileStatement;
StatementTypeOf [WhileSymbol] := ForWithWhileStatement;
StatementTypeOf [WithSymbol] := ForWithWhileStatement;
StatementTypeOf [RepeatSymbol] := RepeatStatement;
end \{ConstantsInitialization\};
procedure WriteA(Character: Char);
var
I: Width;
TestNo: Integer;
$\frac{\text { begin }}{\text { CharCount }}:=$ CharCount +1 ; 01dest $:=$ CharCount mod BufferSize;
$\frac{\text { with UnWritten [01des t] }}{\text { be }}$ do
$\frac{\text { begin }}{\text { if }}$ CharCount $>$ BuffszP1
$\frac{\text { then }}{\text { begin }}$
if ChIsEndLine
$\frac{\text { then }}{\text { begin }}$
$\frac{\text { begin }}{\text { if }}$ IndentAf terEOL $<0$
then
Write(Blanks: - IndentAfterEOL);
OutputCol : $:$ OutputCol - IndentAfterEOL; else $\frac{\text { elise }}{\text { begin }}$
$\frac{\text { begin }}{\text { if }}$ Increment $<0$
if Inc

if $I>0$ then Write (Blanks: I);
TestNo $:=$ LineNumber; $I:=0 ;$
repeat TestNo $:=$ TestNo div $10 ; \quad$ I $:=I+1 ;$
repeat TestNo $:=$ TestNo div 10; I := I
until TestNo $=0 ;$
until TestNo $=0$; $\quad$ ), LineNumber: I);
LineNumber $:=$ LineNumber - Increment;
if LineNumber > 9999
then LineNumber := LineNumber - 10000;
else
$\frac{\text { else }}{\frac{\text { begin }}{\text { WriteLn; }}}$
WriteLn;
if Increment $>0$

```
            \(\frac{\text { then }}{\text { begin }}\)
                        Write (LineNumber: 4, ' ')
                            IineNumber := LineNumber + Increment
                    end
            end:
            then Write(Blanks: IndentAfterEOL) ;
            OutputCol := IndentAfterEOL +1 ;
            ChIsEn
            ChIsEndLIne :=False;
        end \{IF ChIsEndifine
        else.
            begin Write(Ch); OutputCol \(:=\) OutputCol +1
            end \{ELSE\};
    end \(\{\) IF CharGount \(>\) \}; \(;=\) WriteColumn +1 ;
    end \(\{W I T H\}\);
end \(\frac{\text { end }}{\{W r i t e A\} ;}\)
procedure FlushUnwrittenBuffer
\(\frac{\text { begin }}{\text { WriteA( }}\) ();
    with UnWritten:[01dest] do
    begin ChIsEndLine \(:=\) True; IndentAfterEOL \(:=0 ;\) end;
Writécolumn \(:=0 ;\) for \(I:=0\) to BuffSzMI do WriteA('");
```



```
procedure StartNewitineAndIndent;
\(\frac{\text { begin }}{\text { if PackerIsOff and DisplayIsOn }}\)
    then
        begin
            WriteA \(\left({ }^{\prime} \quad \prime\right.\) ); LastSymbol \(:=\) PeriodSymbol
            with UnWritten[01dest!] do
                    begin
            OhIsEndicine \(:=\) True
                \(\therefore\) IndentafterEOL \(:=\) WriteLeftCol + LeftMargin - 1 ;
            \(\therefore\) end;
            end riteGolumn := WriteLeftCol + Lef tMargin;
    end \{IF PackerIsOff\};
    end \{Star tNewLineAndIndent\};
procedure ReadACharacter;
procedure ReadACharacter;
begin ReadColumn \(>\) ReadR ifghtCol.\(~\)
    if R
        begin
            iff ReadRightGol < MaxReadRightcoi
            then begin NextChIsEOL \(:=\) True; ReadIn end
            else ReadCoInumn \(:=2\);
        end
    \(\frac{\text { el.se }}{\text { if }}\) ReadColumn \(=1\) then
            while ReadColưnn \& ReadheftGoi, do
            begin
            if EOITH(Input) then ReadGolumn := 1
            else beginit ReadColumn := ReadColumn +1 ; Get(Input) end
            end;
    if NextChIsEOL \(\quad\) end;
\[
\begin{aligned}
& \text { end }\{\text { IF PackerIsOff\}; } \\
& \text { eStartNewLineAndIndent }
\end{aligned}
\]
```

```
then
    Character := ' '; NextChIsEOL := False; ChIsEOL := True
        ReadColumn := 1;
        if NoFormatting
    then
        WriteA('');
            WriteA('' '); 
            with UnWritten[01dest] do
\[
\frac{\text { begin }}{\text { ChIsEndLine }}:=\text { True; }
\]
            ChIsEndLine := True;
\[
\begin{aligned}
& \text { CIsEndLine }:=\text { True; } \\
& \text { IndentAfterEOL }:=\text { WriteLef tCol }-1 \text {; }
\end{aligned}
\]
            WriteColumn := WriteLeftCol;
        end;
    else
else
    then
        begin
        Character := Input `; ReadColumn := ReadColumn + 1;
        NextChIsEOL := EOLn(Input); Get(Input); ChIsEOL := False;
        NextChIsEOL := EOLn(Input); Get(Input);
```



```
end {ReadACharacter};
        en
```

procedure WriteSymbol;
$\underline{\text { var }}$
$\begin{array}{ll}\text { I: Width; } \\ \text { NumberBlanksToWrite: } & \text { OptionSize; }\end{array}$
$\xrightarrow[\text { begin }]{\text { if DisplayIsOn }}$
then
begin
NumberBlanksToWrite := SymbolGap;
if (LastSymbol in [LeftParenth, LeftBracket, PeriodSymbol]) or
(Symbot irame in inemicolon, RightParenth, RightBracket,
CommaSymbol, PeriodSymbol, ColonSymboll) or (SymbolName in o
CommaSymbol, PeriodSymbol, ColonSymbol]) $\frac{\text { or }}{\text { (SymboliName in }}$ in
[LeftBracket, Leftrarenth]) and (LastSymbol $=$ Identifier)
then NumberBlanksToWrite $:=0$
then NumberBlanksToWrite $:=0$
else（年（SymbolName in AlphaSymbols）and（LastSymbol in AlphaSymbols）

begin WriteA(' '); NumberBlanksToWrite := SymbolGap - 1 ;
end;
主 Writecolumn + Length + NumberBlanksToWrite - 1 >
Writerightcol
then.
Wríteá (");
$\frac{\text { with UnWritten [01dest] }}{6}$ do
$\frac{\text { with UnWritten [01dest] }}{6}$ do
ChIsEndLine := True;
ChIsEnduine $:=$ Irue;
if RackerIsowf
Ehen.
Eken,
Begin InteleftCoil + Lef EMargin + LongLineIndent +
Length-1> Writekightioi

then Length $:=10$ ；
then
－
OptionSize；

IndentAfterEOL $:=$ WriteLeftGol $-1+$ Lef Margin +then
if（LastSymbol in［LeftParenth，LeftBracket，PeriodSymbol］）or （Symbolifame in（Semicolon，RightParenth，RightBracket，
CommaSymbol，PeriodSymbol，ColonSymbol］）or（SymbolName in $\frac{\text { begin }}{\text { end；}}$ WriteA（＇＇）；NumberBlanksToWrite $:=$ SymbolGap－ 1 ； Writecol
if WriteGolumn＋Length＋NumberBlanksToWrite－ $1>$
ightCol．
$\frac{\text { then，}}{\frac{\text { begin }}{\text { Wrí }}}$
Writea（\％）；
biburitten［01dest］do

```
                LongLineTndent;
                TriteGolumn := WriteLeftCol + LeftMargin +
    #M.s:... WriteGoIumin := Writ
```



```
        4.) begin
```



```
    *.* thengmength }:=\mathrm{ WriteRightCo1 - WriteTeftCol + 1;
                TndentsferEOL :=WriteLeftCol - I
                IndentAfterEOL := WriteLeCoL;
            send endy,
        end
        end
        for I := 1 to Iength do WriteA(Symbol[I]);
    end, {IE, DisplawTsOm\;
    LastSymboi := SymboiName;
end {WriteSymbol};
procedure GopyACharacter;
begin
    if DissplayIs0n
    then :
        \mathrm{ begin Write}
            ightCol then
            hi.le (Character ='') and not ChIsEOL do
            not ChIsEOL then StartNewLineAndIndent
            end;
            end;
            then
            begin
            LeftMargin := 0; StartNewLineAndIndent;
            MtMargin := ActualLef EMargin
            end
        else.WriteA(Character);
    end;
end {GopyACharacter}
procedure DoFormatterDirectives
const
                            Invalid = - 1;
type ParamGount = 1..2;
            amGount = 1..2; [
Var Specification: Params;
            FormatOption: Char;
            FormatOption: Char;
        PrevNoFormatting: Boolean
            EndDirectv: CharSet;
    procedure ReadIn(N: RaramCount; var Specification: Params)
    var
. . N
```

```
begin
    fox I := 1 to Nowdo,
        begin
        OopyAdhungitur;
```



```
        then
            Tepeaty*)
            -0ra(óoj);
            OOMC(O
            GopyAGharacter;
            else Speriffication[I] := Invalid;
```


then
begin
Formatotion $:=$ Character
case Fotmat iop tion of
Formatoption:= Character
case romiatoption of

$\frac{\text { begin }}{\text { ReadIn }(1, ~ S p e c i f i c a t i o n) ~ ; ~}$
if (Specification[1] < WriteRightCol - WriteLeftCol - 9)
ox (FormatOption $=$ ' $\mathrm{P}^{\prime}$ )

E':
if Specification $[1]<4$ then
begin
ProcNamesWanted := Specification[1] $>1$;
ProcNamesWanted := Specification[1] $>1 ;$
EndCommentsWanted $:=$ Odd(SPECIFICATON[1]);
end;
$G^{\prime}: \frac{\text { end }}{}:$ SymbolGap $:=$ Specification[1];
G': SymbolGap $:=$ Specification[1];
IndentIndex $:=$ Specification[1];
' ${ }^{\prime}:$ LongLineInde**ecification[1];
'I': IndentIndex $:=$ Specifitation $[1]$
'P': ProcSeparation $:=$ Specification[1];
'S': StatmtSeparation $:==$ Specifical
'S': StatmtSeparation := Specification[1]
end \{CASE\};
end \{SINGLE PARAMETERS\};
begin ${ }^{\text {R }}$
$\frac{\text { begin }}{\text { Read } I n(2, ~ S p e c i f i c a t i o n) ; ~}$
if Specification[2] <> Invalid
then
case Formatoption of
W. (Specification[1] $>0$ ) and (Specification [2] <
$\underset{\text { Buffersize }-2 \text { ) and (Specification [2] }}{\text { (Specification }}$ -
Specification[1] $>8$ )
then
$\frac{\text { begin }}{\text { WriteLeftGol }}:=$ Specification[1];
WriteLeftCol $:=$ Specification[1];
WriteRightGol $:=$ Specification[2];
end;
end; $:=$
f (Specification [1] >0) and (Specification[2]-
' $\mathbb{A}^{\prime}$ : DeclarAlignment := Specification[1];

```
begin \{DoFormatterDirectives\}
    EndDirecty \(\left.:=\left[{ }^{\prime} *^{\prime}{ }^{\prime}{ }^{\prime}\right]^{\prime}\right]\);
```



$$
B^{\prime}, C^{\prime}, D^{\prime}, F^{\prime}:
$$

```
repe
```

```
then
case Formatoption of
if DisplayIsOn
\(\frac{\text { then }}{}\) BunchWanted \(:=\) Character \(='+\) ';
Cr:
begin
PrevDisplay := DisplayIsOn;
Displaylson : = Gharacter \(=\) ' + ';
tif PrevDisplay and not DisplayIson
then
Writea \(\left({ }^{\prime} *^{\prime}\right) ;\) WriteA \(\left.\left({ }^{\prime}\right)^{\prime}\right) ;\)
SayedBunch \(:=\) BunchWanted;
BunchWanted := False;
end
if not PrevDisplay and DisplayIsOn then
\(\frac{\text { begin }}{\text { StáartNemLineAndInđent; WriteA }\left({ }^{\prime}(\prime) \text {; } ; ~\right.}\)
WriteA ('*') ; BunchWanted := SavedBunch; nid \{IF NOT PREV\},
end \(\left\{\right.\) ' \(D^{\prime}\) : \(\}\);
begin
PrevNoFormatting := NoFormatting
Nowrmatting := Chayacter \(=\) '
Displaytison \(:=\) not NoF"ormatting;
if PrevNoFormatting and not NoFormatting
```




```
                end;
end \{Boolean PARAMETERS\}
    end \{GASE\};
els.e \{THEN\} (Giabaciter in Endpirectry) then GopyACharacter;
```




```
end \{DoFormatterDirectives\};
```


## procedure ReadSymbol;

```
cons.t
```

        ReadNextCh \(=\) True;
        DontReadNextCh = False;
    $\underline{\text { var }}$
TestSymbol: Alfa;
CharNumber: Width;
procedure SkipComment;
begin
repeat while Character <> '*' do ReadACharacter; ReadACharacter
until Character $=$ ')';
ReadACharacterf; LastSymbol := Comment; ReadSymbol
end \{SkipComment\};
procedure DoComment;
var
I: OptionSize;
procedure CompilerDirectives;
begin repeat CopyACharacter; until Character in ['[', '*']
end \{CompilerDirectives\};
begin \{DoComment\}
$\frac{\text { begin }}{\text { if }}$ LastSymbol in [Comment, Semicolon] then
begin
LeftMargin :=0; StartNewLineAndIndent;
LeftMargin := Actuallef tMargin;
WriteSymbol; ifi Character $=$ ' $\$$ ' then CompileerDirectives
if Character $=$ ' [' then DoFormatterDirectives;
repeat
whilie Gharacter <> '*' do CopyACharacter; CopyACharacter
$\frac{\text { until. Gharacter }=\text { ')'; }}{\text { CopyACharaeter; LastSymbol }:=\text { Comment; ReadSymbol; }}$
end $\{$ Docomment $\} ;$
procedure ChedikFor(SecondChar: Char; TwoCharSymbol: Symbols;
ReadAllowed: Boolean);
begin
if ReadAllowed then
begin
Length $A=1 ;$ Symbol $[1]:=$ Character;
SymbolName 's= NameOf [Character]; ReadAGharacter;
end;
if Charactert? SecondChar
then beegin
$\frac{b e g i n}{\text { syinderite }}$
Sy inborit 2$]:=$ Character; Length $:=2$;
Symbioiname $:=$
Symboll Name $:=$ TrocharsymboI; ReadACharacter;


$$
\frac{\text { end; }}{\text { d. }}\left\{\begin{array}{l}
\text { CheckFor }\} ;
\end{array}\right.
$$

$$
\text { end }\left\{\begin{array}{l}
\text { end; } \\
\text { \{CheckFor\}; } \\
695 \\
690
\end{array}\right.
$$

then

$$
\frac{\text { case }}{\prime} \text { Character of },
$$

$$
\frac{\text { begin }}{\text { Check }}
$$

$$
\text { egin } \text { CheckFor(\% Comment, ReadNextCh); }
$$

$$
\text { if (Symbolname }=\text { Comment) and PackerIsofi - then } \text { DoComment }
$$

$$
\frac{\text { nI }}{\text { el se }} \text { if Symboliname }=\text { Comment then SkipComment; }
$$

$$
\text { ' } z^{\prime}
$$

## begin

CharNumber := 1; SymbolIsNumber := False;
repeat
Symbol[CharNumber] : := Character; ReadACharacter CharNumber : ChariNumber +1
ntil not (Character in LettersAndDigits)
Length := Charnumber - 1
or CharNumber. $:=$ CharNumber to AlfaLeng do
Pack (Symbo1, I', TestSymbol); I :=1
PascalSymbol[LastPascSymbol] := Tes.tSymbol
while PascalSymbol[I] $\ll$ TestSymbol do $I$ : $:=I+1$
SymbolName := PascSymbo:lName [T].;

begin SymbolisNumber := True; ChaiNumber $:=1$;
repeat
Symbol[ChaiNumber] := Character; ReadACharacter;
CharNumber := CharNumber +1
until not (Character in Digits $+[\cdot \cdot]$ ),
If character in [ ${ }^{\prime} B^{\prime}$; ' $E^{\prime}$ ]
then
begin Symichinvure := Character; ReadAcharacter
CharNumber $:=$ CharNumber +1 ; , $]$
if Character in Digits + ['+', ' $\mathbf{-}^{\prime}$ ] then
Symbol[CharNumber] := Character; ReadACharacter
CharNumber := CharNumber +
until not (Character in Digits)
end;
:= CharNumber - 1 ; SymbolName := Identifier end \{NUMBER\};
begin
repeat ReadACharacter until Character <> ' '; ReadSymbol $>^{\prime}$ end; $:^{\prime}:$ CheckFor $\left({ }^{\prime}={ }^{\prime}\right.$, OtherSymbo1, ReadNextCh $)$; - <'
$\frac{\mathrm{begin}}{\mathrm{Ch}}$
ChecFor ( $=$ =', OtherSymbol, ReadNextCh)
if SymbilName <> OtherSymbol
then CheckFor ( ${ }^{\prime}>{ }^{\prime}$ ', OtherSymbo1, DontReadNextCh) ; end;
if LastSymbol <> EndSymbol

```
thenc(OheckEor('.', Range, ReadNextCh)
    elise SymbodName := PëriodSymbols:
    kegin
        CharNumber := 1
        repeat
            Symbo1[CharNumber] := Character
            GharNumber:% GharNumber + 1; ReadACharacter
            ## [Charumer] := 
            Symbol[CharNumber] := Character;
            *)
            until Character <> '"''; SymbolName := OtherSymbo
            if. Length > WriteRightCoI - WriteLeftGol + 1
        then:
            begin
            TushUwrittenBuffer; Writelm;
            Writeln(' *** STRING TOO LONG.');
            goto }1
        nd end;',}{\mp@subsup{S}{TRING}}{
    end {利 {SE}
    else
    \mathrm{ Seginmoi[1] := Character; SymbolName := NameOf[Character];}
    Length := 1; ReadACharacter
end
end {ReadSymbol}
procedure ChangeMarginTo(NewLeftMargin: Margins);
Var}\mathrm{ IndentedLef tMargin: Margins;
begin
Acgin}\mathrm{ ActualLeftMargin := NewLef tMargin; LeftMargin := NewLeftMargin;
    if LeftMargin < 0 then LefEMargin := 0
    if Lef
        begin}\mathrm{ IndentedLeftMargin := WriteRightCoI - 9-LongLineIndent;
        if Lef tMargin > IndentedLef tMargin
        then LeftMargin := IndentedLef tMargin
end {ChangeMarginTo};
procedure DoDeclarationUntil(EndDeclaration: SymboISet);
procedure DoParentheses
var
SavedLgLnId: OptionSize;
begin
SavedLgLnId := LongLineIndent;
if DeclarAlignment > 0
then
begin
ongLineIndent := WriteColumn + SymbolGap +1 - LeftMargin -
WriteLeftCo1;
epeat WriteSymbol; ReadSymbol
nite Symbolivame \(=\) RightParenth
WriteSymbol; ReadSymbol;
```


## end \{ChangeMarginTo ;

```
procedure DoDeclarationUntil(EndDeclaration: SymbolSet);
procedure DoParentheses
var
```

```
    else
    else.
        LongLineIndent := 1
        ChangeMarginTo(ActualLef tMargin + IndentIndex)
        StartNewLIneAndIndent;
        repeat WriteSymbol; ReadSymbol
        until SymbolName = RightParenth;
        WriteSymbol; ReadSymbol;
    ChangeMarginTo(ActualleftMargin - IndentIndex)
    eng
end {DoParentheses};
procedure DoFieldListuntil(EndFieldList: SymbolSet);
var
            LastEOL: Margins
            AlignColumn: Width;
procedure DoRecord;
    var
        SavedLeftMargin: Width;
    begin
    SavedLeftMargin := Actualluef tMargin; WriteSymbol; ReadSymbol;
    SavedueftMargin :=Actuallef (Hargin; Writesymbol; Reaftym);
    StarENewLineAndIndent; DoFieldListUntil([EndSymbol]);
    ChangeMarginTo(ActualLeftMargin - IndentIndex);
    StartNewLIneAndIndent; WriteSymbol; ReadSymbol;
    ChangeMarginTo(Savedwef tMargin);
end {DoRecord};
procedure DoVariantRecordPart;
    var
        SavedLeftMargin,
        OtherSavedMargin: Margins;
    begin
    OtherSayedMargin := ActualLef tMargin
    If DeclarAlignment > > 0
    then
        \mathrm{ regin , WriteSymbol; ReadSymbol:}
            until SymbolName in [ColonSymbol, OfSymbol]; 874
```



```
            then.
            WrIteSymbol; ReadSymbol;
            with UnWritten[TastEOL] do
                    IndentAfterEOL := IndentAfterEOL + AlignColumn -
                    WriteGolumn; < - , =0
                    if IndentAfterEOL < 0 then IndentAfterEOL := 0;
            end;
            WriteGoIumn := AlignGoIumn
            ChangeMarginfo(ActuaILeftMargin + AlignCoIumn -
            WriteColumn);
            end;
    fif Symb.miName <> OfSymbol then
```


## AlignColumn: Width;

```
procedure DoRecord;
var
```


## SavedLeftMargin: Width

```
egin
GhangeMarginTo(WriteGolumn - \(6+\) IndentIndex - WriteLeftGol) ;
GangeMarginTo (Actuallef tMargin - IndentIndex);
nd \{DoRecord\};
procedure DoVariantRecordPart;
\(\underline{\text { var }}\)
SavedLeftMargin,
OtherSavedMargin: Margins;
\(\frac{\text { begin }}{\text { oth }}\)
Declardiign \(:>0\) Wef Margin
then
xepeat WriteSymbol; ReadSymbol.
If SymbolName \(=\) ColonSymbol
WritteSymbol; ReadSymbol;
with UnWritten [LastEOL] do
ndentafterEOL := IndentAfterEOL + AlignColumn -
if IndeptAfterEOL \(<0\) then IndentAfterEOL := 0
WriteGolumn :=AlignCoIumn;
ginginuo (Actuallef Margin + AlignColumn -
end;
fof Symboiname <> OfSymbol then
```

repeat WriteSymbol; ReadSymbol; until SymbolName $=$ OfSymbol; ChangeMarginTo(Actuailef tMargin + IndentIndex) ;
$\frac{\text { repeat }}{\text { WriteSymboI; ReadSymbol; }}$
If SymbolName <> EndSymbol
$\frac{\text { then }}{\text { begin }}$
StartNewiLneAndIndent;
repeat WriteSymbol; ReadSymbol;
until SymbolName in [LeftParenth, Semicolon, EndSymbol]; if SymbolName $=$ LeftParenth $\frac{\text { then }}{\text { begin }}$

Writ teSymbol; ReadSymbol;
SavedLef $\mathrm{Margin}:=$ ActualLef Margin ;
DoFteldiistuntil ([RightParenth]); WriteSymbo
ReadSymbol; ChangeMarginTo(SavedLeftMargin); $\begin{array}{r}\text { Rea } \\ \text { end; }\end{array}$
end;
until SymbolName <> Semicolon;
ChangeMarginTo(OtherSavedMargin);
end \{DoVariantRecordPart\};

## begin \{DoFieldListUntil

Lasteol := Oldest;
if LastSymbol $=$ Lef tParenth
then for $I:=1$ to DeclarAlignment - Length do WriteA(' ');
AlignColumn $:=$ Lef $\mathrm{Margin}+$ Writeteftcol + DeclarAlignment +1
while not (SymbolName in EndFieldist) do
begin
if LastSymbol in [Semicolon, Comment] then
then begin StartNewlineAndIndent; LastEOL $:=01$ dest end;
then begin SartNewLineAndIndent; LastEOL $:=0$ daes t e
if SymbolName in Recordsymbol, CaseSymbo1, Lef tParenth,
CommaSymbol, ColonSymbol, EqualSymbol]
CommaSymbol, ColonSymbol, EqualSymbol]

## then

case SymbolName of
RecorđSymbol: DoRecord.
CaseSymbol: DotrariantRecordPart;
Lef tparenth: DoParentheses
begin
$\frac{\text { begin }}{\text { WriteSymbol; }}$
if DeclarAlignment $>0$
then
if not (EndLabel $<=$ EndFieldList) $\frac{\text { then }}{\text { begin }}$

With UnFritten[LastEOL] do
Begn AlignColumn - WriteColumn; if IndentafterEOL < 0 then IndentAfterEOL $:=0$; WribeGotumn :=AlignColumn;
if Sisybio 1 Name' $=$ GominaSymbol then:
StartMewilineAndIndent; Lasteol $:=01$ dest end; ${ }^{\text {IF }}$ DeclarAlignment $\} ;$ ReadSymboI;
end $\{,:=\}$
end $\left.\frac{\text { end }\{ }{\{\mathrm{CASE}\}},:=\right\}$

| 892 | else begin WriteSymbol; ReadSymbol end; |
| :---: | :---: |
| 893 |  |
| 894 | end \{DoFieldListUntil\}; |
| 895 |  |
| 896 |  |
| 897 | begin \{DoDeclarationUntil\} |
| 898 |  |
| 899 | StartNewLineAndIndent; WriteSymbol; ChangeMarginTo(Actuallef tMargin + IndentIndex); |
| 900 | StartNewLineAndIndent; ReadSymbol; |
| 901 | DoFieldListUntil(EndDeclaration); StartNewLineAndIndent; |
| 902 | ChangeMarginTo(ActualleftMargin - IndentIndex) ; |
| 903 | end \{DoDeclarationUntil\}; |
| 904 |  |
| 905 |  |
| 906 | procedure DoBlock (BlockName: CommentText; BlockNmLength: Width); |
| 907 |  |
| 908 | $\underline{\mathrm{var}}$ |
| 909 | I: Width; <br> IfThenBunchNeeded: Boolean; |
| 910 |  |
| 911 | AtProcBeginning: Boolean; |
| 912 913 |  |
| 914 | procedure DoProcedures; |
| 915 |  |
| 916 | $\underline{\mathrm{var}}$ |
| 917 | I: 0 . . 20; |
| 918 |  |
| 919 | ProcNmLength: Width; |
| 920 |  |
| 921 | begin |
| 922 | for $I:=2$ to ProcSeparation do StartNewLineAndIndent; |
| 923 | StartNewLineAndIndent; WriteSymbol; ReadSymbol; |
| 924 | for $\mathrm{I}:=0$ to (Length - 1) div AlfaLeng do |
| 925 | Pack (Symbol, I * AlfaLeng + 1, ProcName[I + 1]); |
| 926 | ProcNmLength := Length; WriteSymbol; ReadSymbol; <br> if SymbolName $=$ LeftParenth then |
| 927 |  |
| 928 | begin - |
| 929 | WriteSymbol; |
| 930 | repeat ReadSymbol; WriteSymbol |
| 931 | until SymbolName $=$ RightParenth; |
| 932 | ReadSymbol; |
| 933 | if end Symolvame $=$ ColonSymbol then |
| 934 |  |
| 935 | repeat WriteSymbol; ReadSymbol; until SymbolName = Semicolon; |
| 936 | WriteSymbol; ReadSymbol; <br> ChangeMarginTo(ActualLef tMargin + IndentIndex); |
| 937 |  |
| 938 | StartNewLineAndIndent; LastProgPartWasBody := False; <br> DoBlock (ProcName, ProcNmLength); LastProgPartWasBody := True; |
| 939 |  |
| 940 | ChangeMarginTo(ActualLeftMargin - IndentIndex); WriteSymbol; |
| 941 | ReadSymbol; StartNewLineAndIndent; |
| 942 | end \{DoProcedures\}; |
| 943 |  |
| 944 |  |
| 945 | procedure DoStatement(var AddedBlanks: Width; StatmtSymbol: |
| 946 | CommentText; StmtSymLength: Width) |
| 947 |  |
| 948 | var I: Width; |
| 949 |  |
| 950 | StatmtBeginning: Integer; |
| 951 | StatmtPart: array [1 . 4] of Integer; |
| 952 | B1ksOnCurrntLine, - |
| 953 | BlksAddedByThisStmt: Integer; |
| 954 | Successful: Boolean; |
| 955 |  |
| 956 |  |
| 957 | procedure Bunch(Beginning, Breakpt, Ending: Integer; |

```
StatmtSeparation: OptionSize);
begin
    if BunchWanted or IfThenBunchNeeded
    then
            begin if statmtSeparation \(<1\) then StatmtSeparation :=1;
            BlksOnCurrntLine := BlksOnCurrntLine + StatmtSeparation -
            Successful : \(:\) ( (Ending - Beginning + B1ksOnCurrntLine +
            UnWritten[Beginning mod BufferSize].IndentAfterEOL) <
            WriteRightCol) and (CharCount - Beginning \(<\) BufferSize) ;
Successful
        if Successful
        then
```



```
            IksAddedByThisStmt \(:=\) B
StatmtSeparation -1 ;
            UnWritten[Breakpt mod BufferSize].IndentAfterEOL := -
                StatmtSeparation;
            end;
        end;
end \{Bunch\};
```

procedure WriteComment;
var
I: Width;
SavedLength: Width;
SavedSymbolName: Symbols;
SavedChars: Symbolstring;
begin
SavedSymbolName := SymbolName;
for $I:=1$ to Length do SavedChars [I] $:=$ Symbol[I];
Sor $:=1$ to Length do SavedChars[I] $:=$ Symbol[I];
Symbol[1] $:=\prime$ ( $;$ Symbol[2] $:==^{\prime}{ }^{\prime} ;$ Length $:=2$; WriteSymbo1;
for $I:=0$ to (StmtSymLength - 1) div AlfaLeng do
Unpack (StatmtSymbol[I + 1], Symbol, (I * AlfaLeng +1));
Length $:=$ StmtSymLength; SymbolName $:=$ PeriodSymbol;
LastSymbol:= PeriodSymbol; WriteSymbol; Symbol[1] := '*';
Symbol[2] $:=$ ')'; Length $:=2$; WriteSymbol;
I $:=1$ := Sength do SymboI [I] $:=$ SavedChars [I]
end \{WriteGomment\};
procedure DoStmtList(EndList: Symbols);
var

> B1ksAfterPrt2: Width;
> AtProcEnd: Boolean;
begin
AtProcEnd :=AtProcBeginning; WriteSymbol; ReadSymbol; StatmtPart[1] := CharCount +1 ; StatmtPart[2] := StatmtPart[1]; if SymbolName <> EndList

## $\frac{\text { then }}{\text { begin }}$

if ProcNamesWanted then
if AtProcBeginning then
then if LastSymbol $=$ BeginSymbol then WriteComment;
AtProcBeginning := False,
DoStatement (AddedBlanks, StatmtSymbol, StmtSymLength);
BlksAfterPrt $2:=$ AddedBlanks;
BIksAddedByThisStmt := BlksAddedByThisStmt + AddedBlanks;

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```
            while SymbolName <> EndList do
        begin
            WriteSymbol; ReadSymbol;
        if SymbolName <> EndList
        then
            StatmtPart [3]:= CharCount + 1;
            StatmtPart[3]:= Charcount + 1;
            MoStat ement(AddedBlanks, StatmtSymbo1, 1096
            StmtSymLength);
            B1ksAddedByThisStmt := B1ksAddedByThisStmt +
            Ach(StatmtPart[2], StatmtPart[3], CharCount,
            StatmtSeparation);
            if not Successful
            then
```



```
                    B1ksAfterPrt2 := AddedBlanks;
                end
                else BIksAf terPrt2 := B1ksOnCurrntLine;
        end {WHILE SymbolName <> EndList};
    {IF SymbolName <> EndList};
    lksOnCurrntLine := BIksAddedByThisStmt;
    Bunch (StatmtBeginning, StatmtPart[1], CharCount, SymbolGap);
    SuntNewLineAndIndent; StatmPart[1] := CharCount;
    repeat WriteSymbol; ReadSymbol;
    Til Symbo1Name in [Semicolon, UntilSymbol, EndSymbol,
    E1seSymbol, PeriodSymbol];
    if Successful
    then
            If EndList = UntilSymbol
            then StatmtPart[4] := StatmtSeparation
            else StatmtPart[4]:= SymbolGap;
            Bunch(StatmtBeginning, StatmtPart[1], CharCount,
            StatmtPart[4]);
    end {IF Successful};
    if not (Successful and BunchWanted)
    then
            if LastSymbol = EndSymbol then
            if AtProcEnd and ProcNamesWanted then WriteComment
end {DoS else lftist};
begin {DoStatement}
    BlksOnCurrntLine := 0; Successfu1 := False;
    BlksAddedByThisStmt := 0;
    ChangeMarginTo(ActualLeftMargin + IndentIndex);
    StartNewLineAndIndent; StatmtBeginning := CharCount;
if SymbolIsNumber
then
            with UnWritten[01dest] do
```



```
            if IndentAfterEOL < 0 then IndentAfterEOL := 0;
            end;
            WriteSymbol; ReadSymbol {Write LABEL}; WriteSymbol;
            ReadSymbol {Write COLON};
            ReadSymbol {Write COLON};
    end;
tementTypeOf[SymbolName] of
    ForWIthWhileStatement:
begin \{DoStatement \}
BlksonCurrntLine := 0; Successful := False;
ChangeMarginTo(ActualLeftMargin + IndentIndex);
f Symbinsumer
then
with UnWritten[OIdest] do
\(\frac{\text { Begin }}{\text { IndentAfterEOL }:=}\) IndentAfterEOL - 1 - Length - SymbolGap;
If IndentAfterEOL \(<0\) then IndentAfterEOL :=0;
end;
ForWithWhileStatent:
```

```
begin
    Pack(Symbo1, 1, StatmtSymbol[1]); StmtSymLength := Length;
    repeat WriteSymbol; ReadSymbol
    IName = DoSymbol
    WriteSymbol; ReadSymbol; StatmtPart [1]:= CharCount + 1;
    DoStatement(AddedBlanks,StatmtSymbo1, StmtSymLength);
    BlksOnCurrntLine := B1ksOnCurrntLine + AddedBlanks;
    Bunch(StatmtBeginning, StatmtPart[1], CharCount, SymbolGap);
    end;
RepeatStatement: DoStmtList(UntilSymbol)
IfStatement:
    #egin (Syckmbo1, 1, StatmtSymbol[1]); StmtSymLength := Length;
    Pack(Symbol, 1, StatmtSymbol[1]); StmtSymLength := Length;
        repeat WriteSymbol; ReadSymbol
        StartNewLineAndIndent; StatmtPart[1] := CharGount;
        WriteSymbol; ReadSymbol; StatmtPart[2] := CharCount + 1;
    WriteSymbol; ReadSymbol; StatmtPart [2] := CharGount +
    BlksOnCurrntLine := AddedBlanks;
    BlksAddedByThisStmt := AddedBlanks;
    Bunch (StatmtPa
    then
            Bunch (StatmtBeginning, StatmtPart[1], CharCount,
            StatmtSeparation)
        else IfThenBunchNeeded := True
        else IfThenBunchNeeded := True
        then
            Pack (Symbol, 1, StatmtSymbol[1]);
            Pack(Symbol, 1, StatmtSymboll1]); 
            StmtSymLength := Length; IfThenBunchNeeded := Fais
            StartNewLineAndIndent; StatmtPart[3] := CharCount;
            StatmtPart[4] := CharCount + 1;
            StatmtPart[4] := CharCount + 1; ;
            BlksOnCurrntLine := AddedBlanks;
            BlksOnCurrntLine := AddedBlanks;
            AddedBlanks;
            Bunch (StatmtPart[3], StatmtPart [4], CharCount,
            SymbolGap);
            B1ksOnCurrntLine := B1ksAddedByThisStmt;
            if Successful then
                StatmtSeparation);
        end
        else
        if (Char
            BunchWanted := not BunchWanted;
            B1ksOnCurrntLine := 0;
            Bunch(StatmtBeginning, StatmtPart[1], StatmtPart[2]
                SymbolGap);
            end;
        IfThenBunchNeeded := False;
end {IfStatement};
CaseStatement:
    begin
        repeat WriteSymbol; ReadSymbol
        until SymbolName = OfSymbol;
        WriteSymbol; ReadSymbol;
        ChangeMarginTo(ActualLeftMargin + IndentIndex)
        while SymbolName <> EndSymbol do
    while Sy
            begin
                BunchWanted := not BunchWanted;
hile SymbolName <> EndSymbol do
begin
```

```
```

StartNewLineAndIndent; StatmtPart[1] := CharGount;

```
```

StartNewLineAndIndent; StatmtPart[1] := CharGount;
for I :=0 to (Length - 1) div AlfaLeng do
for I :=0 to (Length - 1) div AlfaLeng do
mmSym
mmSym
until SymbolName = ColonSymbol:
until SymbolName = ColonSymbol:
WriteSymbol; ReadSymbol;
WriteSymbol; ReadSymbol;
if not (SymbolName in [Semicolon, EndSymb.ol])
if not (SymbolName in [Semicolon, EndSymb.ol])
then
then
begin
begin
StatmtPart[2]:= CharCount + 1;
StatmtPart[2]:= CharCount + 1;
DoStatement (Adde CBlanks, StatmtSymbol
DoStatement (Adde CBlanks, StatmtSymbol
tmtSymLength)
tmtSymLength)
BlksAddedByThisStmt := BIksAddedByThisStmt +
BlksAddedByThisStmt := BIksAddedByThisStmt +
AddedBlanks;
AddedBlanks;
Bunch(StatmtPart[1], StatmtPart[2], CharCount,
Bunch(StatmtPart[1], StatmtPart[2], CharCount,
SymbolGap);
SymbolGap);
end {IF NOT(SymboIName...)}
end {IF NOT(SymboIName...)}
if SymbolName = Semicolon
if SymbolName = Semicolon
then begin WriteSymbol; ReadSymbol; end
then begin WriteSymbol; ReadSymbol; end
end;
end;
ChangeMarginTo(ActualLeftMargin - IndentIndex);
ChangeMarginTo(ActualLeftMargin - IndentIndex);
SartNewLineAndIndent; WriteSymbol; ReadSymbol
SartNewLineAndIndent; WriteSymbol; ReadSymbol
if EndCommentsWanted and (LastSymboi = EndSymboI) then
if EndCommentsWanted and (LastSymboi = EndSymboI) then
StatmtSymbol[1] := 'CASE '; StmtSymLength := 4;
StatmtSymbol[1] := 'CASE '; StmtSymLength := 4;
WriteComment;
WriteComment;
end;
end;
end 知; {CaseStatement};
end 知; {CaseStatement};
therStatement:
therStatement:
whil
whil
while not (SymbolName in [Semicolon, UntilSymbol, EndSymbol,
while not (SymbolName in [Semicolon, UntilSymbol, EndSymbol,
begin WriteSymbol; ReadSymbol end
begin WriteSymbol; ReadSymbol end
end {OTHER};
end {OTHER};
CompoundStatement: DoStmtList(EndSymbol)
CompoundStatement: DoStmtList(EndSymbol)
end {CASE};
end {CASE};
AddedBlanks := BlksAddedByThisStmt;
AddedBlanks := BlksAddedByThisStmt;
ChangeMarginTo(ActualleftMargin - IndentIndex);
ChangeMarginTo(ActualleftMargin - IndentIndex);
end {DoStatement};
end {DoStatement};
begin \｛DoBlock\}
LastProgPartWasBody ：＝LastProgPartWasBody and（SymbolName＝ BeginSymbol）；
if SymbolName＝LabelSymbol then DoDeclarationUntil（EndLabel）；
if SymbolName = ConstSymbol then DoDeclarationUntil(EndConst)
if SymbolName = ConstSymbol then DoDeclarationUntil(EndConst)
if SymbolName = VarSymbol then DoDeclarationUntil(EndVar);
if SymbolName = VarSymbol then DoDeclarationUntil(EndVar);
while SymbolName in [FuncSymbol, ProcSymbol] do DoProcedures;
while SymbolName in [FuncSymbol, ProcSymbol] do DoProcedures;
if SymboIName = BeginSymbol
if SymboIName = BeginSymbol
then
then
begin
begin
if Las tProgPartWasBody
if Las tProgPartWasBody
then for I := 2 to ProcSeparation do StartNewLineAndIndent;
then for I := 2 to ProcSeparation do StartNewLineAndIndent;
IfThenBunchNeeded := False; AtProcBeginning := True;
IfThenBunchNeeded := False; AtProcBeginning := True;
DoStatement(T BlockName, BlockNmLength) { I TS DUMMY PARAM };
DoStatement(T BlockName, BlockNmLength) { I TS DUMMY PARAM };
DoStatement(I, BlockName, BlockNmLength) { I IS DUMMY PARAM };
DoStatement(I, BlockName, BlockNmLength) { I IS DUMMY PARAM };
ChangeMarginTo(ActualLeftMargin + IndentIndex);
ChangeMarginTo(ActualLeftMargin + IndentIndex);
end
end
el\mp@code{se begin WriteSymbol; ReadSymbol {Write FORWARD} end}

```
el\mp@code{se begin WriteSymbol; ReadSymbol {Write FORWARD} end}
```

```
            mtSymLength := Length;
```

            mtSymLength := Length;
            repeat WriteSymbol; ReadSymbol
            repeat WriteSymbol; ReadSymbol
            BlksOnCurrntLine := AddedB1anks;
            BlksOnCurrntLine := AddedB1anks;
            ange
            ange
            ElseSymbol]) do
            ElseSymbol]) do
        AdcangeMarginTo BlksAddedByIMIsStmt;
    ```
        AdcangeMarginTo BlksAddedByIMIsStmt;
```

    \(\left.\begin{array}{l}7 \\ 7 \\ 7 \\ \hline\end{array}\right\}\)
    StartNewLineAndIndent;
WriteSymbol; ReadSymbol;
for I :=0 to (Lenth - 1)

MainNmLength $:=$ Length;
repeat WriteSymbol; ReadSymbol; until SymbolName $=$ Semicolon;
WriteSymbol; ReadSymbol; StartNewLineAndIndent;
Droblock (Main, MainNmbength); WriteA('.');
FlushUnwrittenBuffer;
ConstantsInitialization; Initialize;
if EOF(Input) then WriteLn(' *** NO PROGRAM FOUND TO FORMAT.')
ConstantsInitialization; Initialize;
if EOF(Input) then WriteLn(' *** NO PROGRAM FOUND TO FORMAT.')
else
begin \{ CONSTANTS: \}



AlphaSymbols := [ProgSymbol, BeginSymbol, EndSymbol, ConstSymbo
TypeSymbol, RecordSymbol, CaseSymbol, IfSymbo1, ThenSymbol,
TypeSymbo1, RecordSymbol, CaseSymbol, IfSymbo1, ThenSy
ElseSymbo1, DoSymbol, OfSymbol, ForSymbol, WithSymbol,
ElseSymbol, DoSymbol, OfSymbo1, ForSymbol, WithSymbol,
WhileSymbol, RepeatSymbol, UntilSymbol, Identifier, VarSymbol,
ProcSymbol, FuncSymbol, LabelSymbol, Alpha0perator];
EndLabel $:=$ CConstSymbo1, TypeSymbo1, VarSymbo1, ProcSymbol,
FuncSymbol, BeginSymbol];
FuncSymbol,
EndConst $:=$ EndLabel - [ConstSymbol];
EndType := EndConst - [TypeSymbol];
EndVar := EndType - [VarSymbol];
\{ Initialize COLUMM DATA: \}
WriteColumn $:=0 ;$ Lef tMargin $:=0$; Actuallef tMargin $:=0$;
OutputCol $:=1$; ReadLef tCol $:=1$; ReadRightCol $:=$ MaxReadRightCol;
OutputCol $:=1 ;$ ReadLeftCol $:=1$; ReadRightCol $:=$ MaxReadRightCol
WriteLeftCol $:=1$; WriteRightCol $:=$ MaxWriteRightCol; Oldest $:=1$
CharCount $:=1$; LineNumber $:=0$; Increment $:=0$;
\{ Initialize Boolean PARAMETERS: $\}$
PackerIsoff $:=$ True; BunchWanted $:=$ False; DisplayIsOn $:=$ True;
ProcNamesWanted $:=$ True; EndCommentsWanted $:=$ False;
NoFormatting := False;
Noformatting : $=$ False;
\{ Initialize NUMERIC PARAMETERS: \}
IndentIndex $:=3$; LongLineIndent $:=3$; ProcSeparation $:=2$;
$\begin{array}{ll}\text { IndentIndex }:=3 ; \text { LongLineIndent }:=3 ; & \text { ProcSeparation }:=2 ; \\ \text { SymbolGap }:=1 ; & \text { StatmtSeparation }:=3 ;\end{array}$
\{ Initialize INPUT CONTEXT DATA:
ReadColumn $:=1 ;$ ChIsEOL $:=$ False; NextChIsEOL $:=$ False;


end \{Initialize\};
begin \{MainProgram\}
Begin ReadACharacter; ReadSymbol
f SymbolName <> ProgSymbol
then WriteLn(' *** "PROGRAM" EXPECTED.')
else
begin
$* * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * ~$
procedure Initialize；

begi

ReadACharacter；ReadSymbol：
if SymbolName＜＞ProgSymbol
then Writeln（ $* * *$＂PROGRAM＂EXPECTED．＇）
$\frac{\text { else }}{\text { begin }}$

```
I: Width;
\(\xrightarrow{\mathrm{var}}\)
```

$\frac{\text { begin }}{\text { Digits }}:=\left[{ }^{\prime} 0^{\prime}\right.$ ，$\}$

A1phaSymbols $:=$［ProgSymbol，BeginSymbol，EndSymbol，ConstSymbol，
TypeSymbol，Recordsymbol，CaseSymbol，

repeat WriteSymbol；ReadSymbol；until SymbolName＝Semicolon；
Doblock（Main，MainNmLength）；WriteA（ $\left(^{\prime} .^{\prime}\right)$ ；
FlushUnwrittenBuffer；
FlushUnwrittenBuffer；
UnwrittenBuffer；

MOVING A LARGE PASCAL PROGRAM FROM AN LSI-11 TO A CRAY-7 Richard L. Sites, APIS Department, UC/San Diego 92093

In March, 1978, I had occasion to move a 2400-7ine PASCAL program from an LSI-11 at the University of California/San Diego (UCSD) to a Cray-1 at Los Alamos Scientific Laboratory (LASL). At both places, the compiler is a vari-
ant of the P4 portable compiler. This note summarizes the experience and makes several major points about PASCAL:

1. It was possible to move a substantial PASCAL program from a small slow machine to one approximately 150 times bigger. No other language has compatible full-language implementation across such a wide range
of machines--essentially from the world's slowest micro to the world's of machines--essentially
fastest supercomputer.
2. There were compile-time and run-time incompatibilities which should not have existed. The last part of this note is directed to implementors, with a plea to avoid such problems.
3. Using a table-top LSI-11 system, an ori-going project is developing production software for the Cray-1. This would not be feasible in BASIC, FORTRAN, or assembly language
Before describing the problems encountered in moving the program, a little background is needed. The P4 portable PASC.LL compiler is about 4000 lines of $P-C O D E$. P-CODE is the machine language for a pseudo-machine that has a simple stack and about 50 operations. The P-CODE version of a program consists exclusively of a stream of these simple operations, with no associated side tables or assumed information.
On the Cray-1, P-CODE is translated by another 4000 line PASCAL program into Cray-1 assembly language, which then cascades into the standard assembler and oader. This sequence allowed a running, reasonably efficient PASCAL system to be brought up on the Cray-1 with very few months of effort.

On the LSI-11, P-CODE is represented in a very compact form, and is interpreted directly. This has two advantages over compiling to native PDP-11 machine code: First, the P-CODE form of a program is more compact than the machine code, typically by a factor of two. This space compactness is the sole by changing only the interpreter, the identical P-CODE can be run on other micros, allowing the entire compiler and operating system to be transported to cros, allowing

The program moved from UCSD to LASL is the skeleton of a machine-independent optimizer for P-CODE. The initial version of the optimizer will work on Cray-P-CODE, but later versions should work on other variants, and hence one set of optimizing algorithms may eventually be running on a wide variety of machines The characteristics of the two machines and the initial 2400 line skeleton are sumnarized in Table 1.

## Source program moved

```
Pascal source lines 2400
Pascal procedures
P-CODE instructions (Cray-1) 9200
Cray-1 instructions 19100
```

Articles

| Compile times in seconds | LSI-11 | Cray-1 | Approximate ratio |
| :---: | :---: | :---: | :---: |
| Pascal to P-CODE (lines/min.) | 245 (600) | 1.19 (120000) | $200: 1$ |
| P-CODE to Cray-1 asm | - | 2.30 |  |
| Cray-1 asm to binary | - | 4.62 |  |
| Cray-1 loader | - | 0.66 |  |
| TOTALS | 245 | 8.77 (16400) | $30: 1$ |
| Execution times in seconds |  |  |  |
| 75 data lines | 51 | 0.32 | $150: 1$ |
| 2400 data lines | n.a. | 3.43 |  |
| Memory sizes in bytes | 56K | 8000K | 150:1 |

Table 1. Summary of source program moved and machines used.

The rest of this note describes the six major portability problems encountered, along with my suggestions for solutions. Some of these comments parallel those found in other articles on these pages over the past few $1-1 / 2$ days, although I originally expected it to take $1 / 2$ a day. The extra time was wasted on the problems below.

## Portability problem \#1

The Cray-1 compiler recognizes only lower-case ASCII reserved words while the UCSD compiler recognizes only upper-case ASCII. This meant that

This problem clearly subverts the essential idea of ASCII as a standard Code for Information Exchange. It is not sufficient just to have the compiler convert all input to a single case, because (i) character string con stants must not be changed, and (2) ALIAS SPELLINGS of identifiers should not be allowed. An alias spelling is defined to be one that may or may not be recognized as the same as an original spelling, depending on the details declared as:

VAR XYZ : TNTEGER;
could have alias spellings of "xyz" and "Xyz", among others. As a matter of principle, I believe that such spellings should not be allowed because they serve only to introduce confusion about whether the original programmer intended three distinct variables or one. The standard example program for this issue is:
BEGIN
VAR XYZ : INTEGER;
PROCEDURE ABC;
VAR xyz : REAL;
BEGIN
XYZ := 12; (* which block, inter or outer ??*)
END;
Converting all identifiers to upper case resolves the assignment to the REAL variable, while treating the case shifted names as distinct resolves the assign ment to the INTEGER variable. I believe that the original programmer's intent in such a program is truly ambiguous, so the program should not be allowed in the first place. The declaration xyz : REAL should generate a compile-time

## Articles

error (or at least be flagged with a warning) on the basis that an alias spelling of the same variable already exists. Thus, the issue of how to resolve the assignment neve
(1) Reserved words, such as BEGIN are recognized independent of the case of the individual letters, so that "BEGIN", "begin", and "BeGiN" are all recognized as reserved words.
(2) An identifier used in a declaration may have its individual letters in any case, and that particular spelling is inserted EXISTS in the symbol table.
(3) An identifier used in the body of a program must exactly match the spelling in the symbol table, including each letter being of the correct case

These rules allow any program to be compiled, so long as words in it are consistently spelled with the same pattern of upper- and lower-case letters.
(ASIDE: These same rules can be used to detect most cases of identifiers which differ after the first eight letters, without needing to store more than eight letters in the symbol table. Most compiler symbol tables store 7-bit ASCI characters in 8-bit bytes. If all the characters after the first 8 in an identifier are hashed and the hash value stored in the unused bits of these bytes, then rule (2) above can be interpreted to mean "an alias spelling exists (and hence an error/warning message is generated) if some existing identifier in the symbol table has the same first 8 characters, but a different hash code first 8 characters of an identifier and the hash code for the remaining characters must match exactly." This idea completely clears up the concern of A.H.J. Sales (Pascal News, Feb. 1978, p. 78), except when the hash codes for two different tails turn out to be identical; this can be made rare, and can be guaranteed not to happen for single-character differences. End of ASIDE.)

## Portability problem \#2

Contrary to the Report, the Cray-1 compiler does not recognize empty field ists in variant record declarations, RECORD CASE I: BOOLEAN OF TRUE: (X,Y INTEGER); FALSE: ()END, nor does it recognize untagged variant records, RECORD CASE BOOLEAN OF .... The lesson here is clear--recognize the entire language as defined, without taking shortcuts.

## Portability problem \#3

UCSD Pascal includes non-standard procedures OPEN and CLOSE. I had to rework the calls to use the standard RESET and REWRITE, which lack two useful capabilities: (1) there is no way to close a file explicitly, and hence there is no way to release a file for other uses before the program terminates; ther disk file) or kept (a disk output file) after termination; (2) there is no way to open a file explicit.ly, supplying a character-string file name at that time These are limitations I can live with, but I would prefer to see some agreedupon standard extensions in this area

## Portability problem \#4

## TYPE WHOLENUM $=0 . .32767$ <br> VAR I : INTEGER;

FUNCTION F(...):WHOLENUM;

## I : $\because \dot{F}(\ldots) * I$;

while the Cray-1 compiler complained about operand incompatability at the multiply. Inconsistent type checking is a well-known problem in Pascal, so I won't dwell on it. In this particular case, though, I am frustrated because the whole purpose of introducing the type WHOLENUM is to convey to the reader (and the compiler) the idea that all WHOLENUMS are intended to be non-negative My temporary fix was : TYPE WHOLENUM = INTEGER;

At this point in the process, my 2400 Tine program compiled properly and executed for the first time. I had fixed problem \#1 by converting the entir executed for the first time.
program to lower case, and this fix now came back to haunt me, because the input data file was still in upper case, and hence did not match any of my lower-case character-string constants. Converting the entire data file to lower case also did not quite do the trick, because my program's output (remember, the 2400 ine program optimizes P-CODE) cascades eventualiy into the Cray-1 loader, which demands standard procedure names (such as SIN) in upper case. Clearly, the case shift problem was taking more energy than it should.

## Portability problem \#5

The first real problem to crop up in execution was that my hash function always returned the same value, zero, instead of reasonably we ll distributed series of numbers in the range 0..127. The hash function was built using (conceptual) shifts and exclusive-or's, and in fact did a fair amount of lying with variant records to jump between character, integer (I+I used for left shift of one bit), and set (S1+S2, S1*S2, and ALLBITS-S used to build XOR) representations. There is a serious issue here of how to build a portable remove variant records as an "escape hatch". Try it yourself -- build a function which accepts a PACKED ARRAY [0..7] OF CHAR and returns an integer in the range $0 . .127$. The particular hash function desired XORs the 8 characters, each one offset one bit from the next to get a 5 -bit intermediate. The upper 8 bits and the lower 7 are then XORed, and the lower 7 bits of this are returned. (This particular function guarantees different hash values for in puts which differ by any one character, or which differ by transposition In addition, no overflow is generated on a l6-bit machine.)

## Portability problem \#6

The final output of the program was spaced funny. In the statement WRITELN $(3,4)$;

UCSD Pascal inserts no blanks around the fields, giving "34", while Cray-7 Pascal uses a default field width, giving " 3 4". In my application, the blanks are not wanted, but Pascal output editing is not precisely defined, so many implementations supply extra blanks. Often, these blanks reflect a legitimate desire to separate items of output when specified field Widths are exceeded, as in WRITELN(100:2, 200:2) which normally prints as " 100 200". I propose that a specified field width of zero mean no padding

Overall, moving a 2400 -line Pascal program froved surprisingly successful,
having done it once should make it easy to move a 5000 -line program this summer.

On the Article "What to do After a While"
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## INTRODUCTION

The letter by A.H.J. Sale ${ }^{1}$ and the article by Barron and Mullins ${ }^{2}$ in PASCAL News \#11 address themselves to an ambiguity in the definition of the PASCAL language: should Boolean expressions be evaluated in a parallel or sequential manner?

For example, when we write " $P$ and $Q$ ", do we mean

1) (parallel or "logical" evaluation)

$$
P \wedge Q \quad(=Q \wedge P)
$$

or
2) (sequential evaluation)

$$
\text { if } P \text { then } Q \text { else false }
$$

I argue here for the parallel approach.

## THE PROBLEM OF PSYCHOLOGICAL SET

My first objection to sequential evaluation is that it looks parallel to anyone who has had any exposure to symbolic logic. This is the problem of "psychological set", first discussed by Gerald M. Weinbers ${ }^{3}$. This term connotes a state in which our way of thinking about a situation blinds us to its reality. For example, a common error encountered by programmers who use languages in which variables need not be declared is the use of misspelled variable names which "look like" other (valid) names. This kind of error can be extremely hard to find.

Thus, a maintenance programmer who runs into the expression

$$
\text { while }(i<=\operatorname{maxsize}) \text { and }(a[i]<>i t e m) \text { do }
$$

and later finds
while (a[i] <> item) and (i<=maxsize) do may not even see them as different expressions!

## PROVING PROGRAMS CORRECT

The advent of structured control statements has generated a great deal of interest in the problem of proving, either by hand or automatically, the correctness of programs ${ }^{4}$. My second objection to symbolic evaluation is that it will probably in-
crease the difficulty of doing such verification by an order of magnitude. This belief is based on the fact that, in abstract mathematics and logic, non-commutative (i.e., ordexdependent) objects are much harder to handle than commutative objects.

## ON "THE SPIRIT OF PASCAL"

Barron and Mullins argue that sequential evaluation allows us to program "more in the spirit of PASCAL". Whatever that patriotic remark means, I strongly disagree. Let's look at the example they give. We are to search a table for a given item. Using sequential evaluation, their solution is:
var table : array [1..maxsize] of whatever;
-
index := 1;
while (index <= maxsize) and (table[index] <> item)
do index := index + 1;
(* condition for item not found is "index > maxsize" *)

## There are two fundamental flaws in this solution:

1) The solution trlists the algorithm to fit a given data structure.

One of the great advantages of PASCAL over most other languages is the ability it gives the user to create data structures which work well with a given problem. Consider the following solution to the table search problem ${ }^{5}$ :
var table : array [0..maxsize] of whatever;
-
table[0] := item; (* put in sentinel for end of search *) index := maxsize;
while table [index] $<>$ item
do index := index - 1 ;
(* condition for item not found is "index $=0$ " *)
2) The repetitive construct mixes together logical and iterative repetition.
Indeed, in Algol $68^{6}$, a cleaner way to write the BarronMullins algorithm would be:
index := 1;
for $i$ from 1 to maxsize while table $[i]<>$ item
do index := i + 1 od;
\# condition for item not found is "index > maxsize" \#

The problem with this technique is that on exiting the loop, one does not know if termination was caused by the count being exceeded or by the logical condition failing. This is a common error-causing situation, better known as "exiting a loop to the same place from the side and the bottom" ${ }^{7}$.

## ON "EFFICIENT" ALGORITHMS

Finally, $I$ would like to take a moment to talk about efficiency. Barron and Mullins say:
"... But in the Pascal community we should have gotten beyond judging features solely in terms of implementation efficiency. What matters is being able to write correct programs that are easily comprehensible."

The answer to the question of efficiency is not so simple. True, the first consideration of the designer should always be the correctness and clarity of the design. But efficiency often comes in a close second, and sometimes it's a dead heat: in certain circumstances (e.g., real time applications or CAI) if the program is not efficient enough, in terms of either size or execution time, it is irrelevant as to whether or not it's clear, or even correct - the program is unusable.

As Donald Knuth ${ }^{8}$ and others ${ }^{9}$ have pointed out, the problem is generally not that the designer has made efficiency a consideration, but how he has tried to make the design efficient. The villain is not efficiency itself, but micro and premature optimization.

In fact, the algorithm given in this article is a good example of how to optimize a program: by improving its data structures and algorithms. In a Ratfor preprocessor written in PASCAL, the substitution of the algorithm given above for the original one (which was essentially the Barron-Mullins algorithm) resulted in a $30 \%$ decrease in the preprocessor's execution time with no impairment of the clarity of the program.

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(* Receiyed 78/05/11 *)

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A Resolution of the boolean expression-EVALUATION QUestion
or
IF NOT PARTIAL EVALUATION THEN CONDITIONAL EXPRESSIONS

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Introduction
The programming languages ALGOL-60[1] and ALGOL-W[2], which contain the precursors of many of the elegant features of PASCAL, are richer than PASCAL in the variety of ways that an expression may be formed. Both ALGOL-60 and ALGOL-W contain the conditional expression and ALGOL-W contains, in addition, the case expression and the value block. A PASCALized" summary of these constructs is shown in the syntax diagrams below.
relational expression

expression


Although the effects of the conditional expression, case expression, and value block may be had in PASCAL (or in FORTRAN, for that matter), the resulting constructs require multiple statements and the declaration and use of temporary variables that are not otherwise incorporation of these forms of expression into PASCAL on the following grounds:

1. The increased programing facility that they offer more than
compensates for the increased syntatic complexity which their adoption would entail. 2. The conditional expression, in particular, promotes rigor by
removing semantic ambiguities that exist in the evaluation of Boolean expressions. 3 . None of these extensions conflicts with the PASCAL design goals cited by Vavra[3].

We shall restrict the scope of this paper to the case for the conditional expression, showing first some examples of its use, the and second the way in which it avoids the current arguments concerning the proper evaluation of Boolean expressions

The Conditional Expression
In BNF notation, the <conditional expression> may be defined to be
if <Boolean expression> then <expression> else <expression>
(ALGOL-60 restricts the <expression> following the then to be a <simple expression>.)

This construct permits such statements as:
a. $x:=($ if $n<100$ then $a+b$ else $a-b)$ * ( $c+d)$;
b. $a:=\operatorname{sqr}(b[$ if $i$ in $s$ then i else 0]);
c. while if $x<=10$ then $a[x]<>b$ else false do $x:=x+1$;
d. append $:=\frac{\text { if }}{x}$ null ( $x$ ) then copy ( $y$ )
$\frac{\text { else }}{\text { cons }} \overline{x^{\wedge}}$.class $\left(x^{\wedge}\right.$.car) apt then
else referenceerror('append: ist arg invalid structure');
Expressing these constructs in PASCAL is straightforward. Example c could be written as

$$
\begin{aligned}
& \text { found : }=\text { false; } \\
& \text { while }(x<=10) \text { and not found do } \\
& \text { if } a[x]=b \text { then found:=true else } x:=x+1
\end{aligned}
$$

Thus, the PASCAL version of $c$ requires two statements and the extra
Boolean variable, "found". Boolean variable, "found".

Resolution of a Semantic Problem
Another possibility is that the previous example could be written as

$$
\text { while }(x<=10) \text { and }(a[x]<>b) \text { do } x:=x+1 \text {; }
$$

provided the evaluation of the Boolean expression is terminated as soon as $x<=10$ becomes false. This avoids errors when a[ll] does not exist or when it is undefined. This approach is currently the subject of some
debate. In two recent articles in Pascal News, the authors Sale[4] and Barron and Mullins[5] have taken opposite positions regarding standards for the evaluation of Boolean expressions. Sale recommends the "Boolean operator" approach which forces full evaluation of the complete expression, whereas Earron and Mullins prefer "sequential conjunction" which permits the compiler to terminate evaluation of an expression as
soon as its truth or falsity is unequivocally determined.

The reasons that have been given for partial evaluation seem to be

1. efficiency 2. the resolution of cases in which one or more of the terms and factors of the expression are undefined.

Although Barron and Mullins have described three syntactically correct ways of avoiding the problem of point 2 by segmenting the expression, still they advocate the use of partial evaluation. Their
position is understandable, for the techniques required are contrived. Unfortunately, they are the only reasonable ones available with the present language, and partial evaluation makes the code appear to be simpler.

The User Manual [6], as noted by Sale, interprets the Report[6] neither to require nor to forbid the full evaluation of Boolean expressions. However, the syntax of the <expression> given in the Report clearly implies that all operators are to be applied in the evaluation of an expression. Thus, it seems reasonable to expect that any action which appears explicitly in the flow of control must b evaluated. If this is not the case there will always be an uncertainty as to what portions of the program have been executed. For example, the statement

## while A and B do ... ;

means that the statement following do is to be executed if $A$ and $B$ are both true. According to Barron and Mullins, this would be reinterpreted to mean "don't evaluate $B$ and don't execute if $A$ is false." There is a subtle difference between these two notions. The difference is important because $B$ might necessary operations on global variables or var parameters.

Full evaluation of the expression is in keeping with the syntax described in the Report and with intuition. From the language-design standpoint, there seems to be no justification for performing a partial evaluation of an expression. This is particularly true since the most compelling reason advanced for the partial evaluation is to avoid an awkward temporary variable. The conditional expression is a complete the programmer, of the terms and/or factors that are to be evaluated

Example c, above, solves problem 2 by explicitly directing the flow of control around impossible cases. It does not depend on implicit conventions of partially evaluating expressions.

In our opinion, the only reason for not fully evaluating an expression is efficiency of time and memory utilization. While the use
of partial evaluation does have an advantage over the standard PASCAL of partial evaluation does have an advantage over the standard PASCAL construction, the advantage is insignificant when it is compared against
the conditional-expression approach. The following shows the code segments that might be generated for example $c$, above, if the target machine wese a PDP-11.

Conditional Expression
s: evaluate the condition put result of $x<=10$ on stack

```
CMP (SP)+,#TRUE
```

BEQ $2 \$$

> evaluate the else expression put false on the stack

BR

2\$: evaluate the then expression put $a[x]<>b$ on stack
$\begin{array}{ll}3 \$: & \text { CMP } \\ & \text { BNE } \\ & 4 \$\end{array}$
perform do
BR $1 \$$
4 \$:

Partial Evaluation
1\$: evaluate expression put $x<=10$ on stack

CMP (SP),$+ \# T R U E$ BNE $4 \$$
evaluate expression put $a[x]<>b$ on stack
CMP (SP)+,\#TRUE
BNE 4 S perform do

BR l\$
4 :
We see from these examples that the space advantage gained from partial evaluation is that for evaluating the else part of the if and a branch instruction. In this case it amounts to two instructions. It is more the else condition is not evaluated. We feel that there is insufficient justification for adopting partial evaluation as a standard feature of the language. It might, however, be a desirable implementation-dependent feature activated by a compiler directive.

## Conclusions

PASCAL is not yet a complete language in that inclusion of several desirable features of other languages has not yet been openly debated. We recommend the case expression, the value block, and particularly the conditional expression as additions to the language. The basis for this recommendation is that these features will promote semantic rigor, will not conflict with any language-design goals, will provide the programmer with new and useful tools, and will improve the efficiency of the generated code over standard PASCAL

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(* Received 78/08/07 *)

## What to do after a while .. longer

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| case 1234 |  |  |
| :---: | :---: | :---: |
| $\mathrm{p}=\cdot$ $\mathrm{q}=$. $=1$ | expression | Boolean operator expression |
|  | $\begin{aligned} & \text { false } \\ & p \text { and } q \\ & p>q \\ & p \end{aligned}$ | not ( not p or not q ) $p$ and not $q$ |
| -1 - | $\mathrm{p}<\mathrm{q}$ | not $p$ and $q$ |
| $\cdot$ . . . $1 \begin{aligned} & 1 \\ & 1\end{aligned}$ | $\begin{aligned} & q \\ & p<>q \\ & p \text { or } q \end{aligned}$ | not $p$ and $q$ or $p$ and not $q$ not ( not $p$ and not q) |
| $1 \cdot$ • 1 | not ( p or q $\mathrm{q}^{\text {) }}$ | not $p$ and not $q$ |
| $1 \cdot 1 \cdot 1$ | pot ${ }_{\text {not }}^{\text {q }}$ | $p$ and $q$ or not $p$ and not $q$ |
| 1.11 | $\mathrm{p}>=\mathrm{q}$ | p or not q |
| 11. | not p |  |
| 11 ; 1 | p<=q ${ }^{\text {not }}$ ( $p$ and $\left.q\right)$ | $\begin{aligned} & \text { not } p \text { or } q \\ & \text { not } p \text { or not } q \end{aligned}$ |
| 1111 | true |  |

If we were starting the language design again and
We wanted to include a facility for telling implementations how to evaluate expressions (though, in view of B\&M's own so that we were looking for "sequential conjunction" versions of
then I, for one, would oppose the use of "and" \& "or" themselves, on the ground of their old and strong, Boolean algebra connotation.
If POP-2 and RTL/2 have already "adopted" (it should be "adapted") them for such a purpose, that is their problem.
Not that I care, but B\&M's function andop seems unnecessarily complicated to me
function andop ( $p, q$ : Boolean) : Boolean ;
is sufficient because the arguments $\frac{\text { end }}{a r e}$; both evaluated when the function is called - surely ?

What the spirit of Pascal says to me is that we ought not to (i) write programs that rely on not well-defined factors side-effects of functions or undefined values,
(ii) depend on implementors to let us get away with them,
(iii) tell implementors to let us get away with them,
or (iv) complain if implementors use any means they can
or (iv) complain if implementors use any means they can devise to prevent us getting away with them.

The spirit of Pascal also says that it rather fancies itself as a two-edged sword!

$$
78 \text { se } 20 \mathrm{~W} \quad \text { Mintuat Imith }
$$

$+-* /:=., \quad ; \quad 1=\langle \rangle\left\langle<=\gg=()\left[. \quad\right.\right.$ 〕 $\uparrow \ldots\left({ }^{*}{ }^{*}\right)$
copies to :- D.W.Barron J.M.Bishop K.Jensen G.H.Richmond A.H.J.Sale N.Wirth

* Received 78/09/26 *)


## Laurence V. Atkinson

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## In a nutshe 11

A number of recent articles have highlighted problems with multiple exit loops in Pascal. Many of these problems disappear when a loop is controlled by a user-defined scalar

## Introduction

Multiple exit loops and problems with their implementation have featured prominently in four reaent articies: Barron and Mullins [2], Bishop [3], Bishop 4. and Hoxton [5]. Many of these problems do not occur if user-defined scalar constitutes a multi-state process. Pascal's ordinal types provide a natural means of identifying multiple states. This state transition approach is introduced by first considering the Barron and Mullins paper [2] and then taking the other articles in turn.

## Barron and. Mullins

Their example is linear search for a specified item within a vector (assume fuli) but considering the possibility that the desired item may be absent, The program they produce is
const maxsize $=\ldots ;$ succmaxsize $=\ldots ;$
van table : array $[1 .$. maxsize $]$ of whatever;
index : l...succmaxsize;
index := 1 ;
while (index $<=$ maxsize) and (table [index] <> item)
do index := index +1
if index $>$ maxsize then $\{$ item absent\} . . . else $\{$ item found\} . . .
Barron and Mullins claim that "this is a natural way of expressing the peration to be carried out" and is inkeeping with "the spirit of Pascal" The point of their paper is that this program is viable only if boolean basis. The Report [6] leaves this issue open but the User Manual [6] states that all operands in a boolean expression will be evaluated. Jensen and wirt [6] (Chapter 10) produce an equivalent example to illustrate the problem. A state transition approach to their solutions is presented by the present author in [1].

I suggest that a programming style both more natural and more in the spirit f Pascal is achieved when user-defined scalars are introduced and used as state indicators.

State Indicators
In a simple search environment there are three distinct states of interest:
(i) I haven't found it yet but I'm still looking,
(ii) got it,
(iii) I've looked everywhere but it's not here.

This leads us to a solution using a three-state scalar (figure 1)

## const endoftable $=\ldots$;

type toendoftable $=1 \ldots$ endoftable;
searchstates $=$ (searching, thingぁsient, thingfound);
var item : array [toendoftable] of things;
here : toendoftable;
outcome : searchstates;
here $:=1$; outcome $:=$ searching
repeat
if item [here] $=$ thingwanted then outcome $:=$ thingfound else
if here $=$ endoftable then outcome $:=$ thingabsent else here : $x$ succ (here)
until outcome <> searching;
case outcome of
thingfound : . . ;
thingasent :. .
end \{ease\}
Figure '1. Linear search with state transition

We now comment upon the program of figure 1
(i) The intent of the program is more readily apparent.
(ii) The program is now more easily extended to include other cases of
(iii) Subsequent processing, upon exit from the loop, is more transparent:determination of whether or not the desired item has been located is cleaner.
(iv) The compound boolean expression has disappeared and so issues of
boolean operator' or 'sequential conjunction' approach are avoided.
(v) The order of making the tests is not implementation dependent:-
the desired order is unambiguously expressed.
(vi) The subscript cannot go out of the bounds of the array. In Barron and Mullins version the range of the subscript must be one greater than the index range of the array. We return to this point when discussing Bishop's paper [3].
(vii) No redundant tests are made. Barron and Mullins incur a test (index $\varsigma=$ maxsize) which is'always true upon entry to the while loop. My objection to redundant testing is based on considerations of logic rather than efficiency.

```
searchstates = (searching, absent, foundinfirsthalf, foundinsecondhalf);
repeat
    if item[here] = thingwanted then
        case here <='(endoftable div 2) of
            true : outcome := foundinfirsthalf;
            false : outcome := foundinsecondhalf
        end {case}
    else . . .
until outcome <> searching;
case: outcome of
    foundinfirsthalf : . . . ;
    foundinsecondhalf : . . .
    absent : . . .
end {case}
Figure 2. Extended linear search with state transition.
```


## Rishop

Judy Bishop [3] addresses the general problem of subrange exhaustion in a loop of the form

$$
\begin{aligned}
& \text { i : =min; } \\
& \frac{\text { while }}{\text { begin }}(\dot{i}<\max ) \text { and condition do } \\
& \frac{\text { begin }}{\text { son }} \\
& \text { \{ : Something }{ }^{=} \text {succ }(i) \\
& \text { end }
\end{aligned}
$$

in conjuction with

$$
\begin{aligned}
& \text { type } \text { index }=\min \ldots \text { max; } \\
& \underline{\text { var }} i: \text { index; }
\end{aligned}
$$

When considering Barron and Mullins it was noted that one natural consequence of the state transition approach was that the subrange variable could not exceed its bounds. The present problem is therefore solved by this same approach and as before, produces a more transparent program (figure 3). The point raised by John Strait, and discussed by Judy Bishop in [4] is also covered by this approach. HCEtRR

Mark Horton [5] considers two examples each involving a double exit loop and uses them as a basis for suggesting a modification to the Pascal language. He encourages the use of a deterministic loop which, without any indication
type index $=$ min.. $\max ;$
var i : index;
state : (looping, rangeexhausted, otherexitcondition)
i := min; state $:=$ looping
repeat
\{something\}
if $i=\max$ then state $:=$ rangeexhausted else
if $\cdots$ then state $:=$ otherexitcondition else

$$
i:=\operatorname{succ}(i)
$$

until state $\langle>$ looping;
case state of
rangeexhausted : . . . ;
otherexitcondition : . . .
end \{case\}
Figure 3. Bishop's loop with state transition.
of the fact at the loop control level, can jump completely out of itself and far away I do not claim that state indicators can remove the need for all gotos but they can We conider them in turn.

1. Binary search

Horton's program is

$$
\begin{aligned}
& \text { const maxsize }=\ldots \text {; succmaxsize }=\ldots \\
& \text { var a : harray [1...maxsize] of } \cdots \text {; } \\
& \ell: 0 . . \text { maxsize; } \\
& \text { found : boolean. } \\
& \ell:=1 \text {; u := n; } \\
& \text { loop while }(\ell<=u) \text { flag found do } \\
& \text { mid := ( } \ell+\mathrm{u}) \text { div } 2 \text {; } \\
& \text { if } x \leqslant a[m i d] \text { then } u:=m i d-1 \text { else } \\
& \text { if } x>a[\text { mid] then } l:=\text { mid }+1 \text { else } \\
& \text { exit found } \\
& \text { end; }
\end{aligned}
$$

which is a sỵntactic sugaring of the following true Pascal fragment.

R: Labey 1 ;


- witlie $\ell<u$ do
$\frac{\text { begin }}{\text { mid }}$
mid $:=(l+u)$ div 2 .

$$
\text { eng } \quad \cdots ;
$$

1 : 皆d; found then $\cdots$ else ..
Again we find our familiar three-state process.
Although we should not worry unduly about minor points of efficiency we must still bear overall efficiency considerations in mind when designing an algorithm. To be most aesthetically pleasing one of the first tests a program should make in a search loop is 'is what I'm looking at what I want?'. However, for binary search, we suffer if we test for equality before we test relative magnitude. This is becatise, in general, we will hit elements we don't want far more often than we hit an element we do want. Consequently, for about half of our probes, we shoul know which pointer to move after making only one comparison. Accordingly we program.
(i) Program intent is more transparent:
(ii) No modification to the language is necessary to permit a clean solution
(iii) The subscripts cannot go out of the bounds of the array. Horton's program suffers from a variant of Bishop's problem: if the sought entry falls outside the table Horton's version terminates with $\ell-\mathrm{u}=1$
(iv) The ner progran is more
interest. In particular we may be interested to know if we found an item on the final probe available (ie when (top $=$ bottom) and (itemwanted $=$ itemat [top]) \} or earlier (in which case itemwanted $=$ itemat [middie]).

The computation in both programs is the same but for the extra test
'state <> stillchopping' now at the end of each iteration. This test can be implemented (by any compiler anticipating this form of loop control) as a single jump (jump on zero) so this overhead should be of little concern to us.
2. Prime numbers

Horton's program is
const $n=\cdots ;$
var $p, d: 2 \ldots n$;
potential_prime : boolean;
loop for $\mathrm{p}:=2$ to n flag potential_prime do
$\frac{\text { lqop for }}{\text { if }} \mathrm{d}:=2$ to trunc (sqrt d ) ) $\mathrm{do}=0$ then next potential prime
$\frac{\text { end; }}{\text { write }}(p)$
end

$$
\begin{aligned}
& \text { if } x<a[m i d] \text { then } u:=\text { mid }-1 \text { else } \\
& \text { if } x>\text { ath } \\
& \text { Eound := trues goto } 1
\end{aligned}
$$

const endoftable $=\ldots$;
type span $=1$.. entototabie
var itemat : array [span] of ...
bottiom, middle, top : span;
state : (stillchopping, found, absent);
bottom := 1; top := endoftable; state := stillchopping;
repeat
if top $=$ bottom then
qase itemat [top] $=$ itemwanted of
true \& state $:=$ found;
false : state := absent

## end \{case\}

else
begin
middle := (top + bottom) div 2
if itemwanted < itemat [midd1e]
then top $:=$ middle -1 else
if itemwanted $>$ itemat [middle]
then bottom := middle +1 else
state := found
end
until state <> stillchopping
case state of
found : . . . ;
absent :

Figure 4. Binary search with state transition
which, without the syntactic sugar, is
label 1 ;
bonst $n=\ldots$;
var $\mathrm{p}, \mathrm{d}: 2 \ldots \mathrm{n}$;
potentialprime : boolean;
potentialprime := false;
for $p:=2$ to $n$ do
Eegin for $:=2$ to trunc $(\operatorname{sqrt}(p))$ do
if $p, \frac{\text { mod }}{} \mathrm{d}=0$ then
$\frac{\text { begin }}{\text { potentialprime }:=\text { true; }}$
ite $\frac{\mathrm{end} \text {; }}{(\mathrm{p}) ;}$
1 : end

Horton mentions that only odd numbers and divisors need be tested. In the finite state approach we still sweep through contiguous numbers (although we could avoid it - as could Horton) but this time start at 5 and test only odd divisors Since divisors start at 3 it is sensible to make the loop deal with primes $>3$ (hence $>=5$ ). Accordingly primes $<=3$ are best dealt with separately. Apart
from these modifications we stick to Horton's algorithm (figure 5). There should be no need to reiterate previous comments.

```
const \(\mathrm{n}=\ldots\);
var \(p\), potfactor, rootofp : 2 .. \(n\);
    state : (moredivisors, factorfound, pisprime)
if \(n \ll 3\) then primesupto ( \(n\) ) else
begin
    primesupto (3);
    for \(p:=5\) to \(n\) do
        if \(p\) mod \(2<>0\) then
            begin \{ \(p\) is odd \}
                rootofp := trune(sqrt:(p))
            potfactor \(:=3\); state \(:=\) moredivisors;
            repeat
                    if \(P\) mod potfactor \(=0\) then state \(:=\) factorfound else
                    if potfactor \(\rangle=\) rootofp then state \(:=\) pisprime ellse
                    potfactor := potfactor +2
            until state \(\langle>\) moredivisors;
            if state \(=\) pisprime then write ( p )
            end \(\{p\) is odd \(\}\)
```



```
                            Figutate 5! Mrime ntambers' with state transition.
Fightie' 5 !
-
```


## conclustidr

We have seen some illustrations of a particular style of programming The state transition technique is applicable to a number of programming situations and to multi-exit loops in particular. I (and my students) have adopted this approach for a number of years and have rarely suffered from Barron's, Horton's or Mullins/Bishop's complaints. My response to Barron and Mullins! query "What to do after a while?" is "Know the state you are in!

Referdrins
[1] L.V. Atkinson, "Pascal scalars as state indicators", 1978 (under review)
[2] D.W. Barron and J.M. Mullins, "What to do after a while", Pascal News, \#11, 48-50, 1978
[3] Judy M.•Bishop, "Subranges and conditional loops", Pascal News \#12, 37-38, 1978.
[4] Judy M. Bishop, Letter to John Strait, Pascal News, \#12, p51, 1978.
[5] Mark D. Horton, Letter to the editor, Pascal News, \#12, 48-50, 1978.
[6] Kathleen Jensen and Niklaus Wirth, Pascal - User Manual and Report, Springer-Verlag, 1978.
(* Received 78/09/15 *)

# Open Forum for Members 

# DEPARTNENT OF DEFENSE DEPENDENTS SCHOOLS <br> DARMSTADT CAREER CENTER 

 APO Niew York 09175SUBJECT: Pascall News

Mr. Andy Mickel
University Computer Center
227 Experimental Engineering Minneapolis, MN 55455

June 8, 1978

TO: Pasceal User's Group
c/o Andy Mickel
University Computer Center 227 Ex
208 S.E. Union Street
University of Minnesota
Minneapolis, Minnesota 55455

Dear Andy,
Oum school computer group reads with great interest the developments you have presented in Pascal News and additional papers obtained from UCSD written by Kenneth L. Bowles concerning Micro Computer Based Mass Education and the Personalized System of Instruction (PSI). Untill such time when we can pilot and implement microbased systems, we would like, as a first step, to obtain a Pascal implementation for $7 / 16$ 's and 3 Univac $90 / 30$ 's with some 200 terminals supponting BASIC. This office represents the European region and has special interest in the Interdata 7/16 implementation. (Our Pacific region operates the Univac $90 / 30$ 's.) These systems are devoted to our instructional program where one of our goals is to generate computer literate youth.
Many of our graduates become employed with DOD in some capacity.
DOD, with its "Ironman" project is moving towand a Pascal based If standard system, it is only natural for its own school system to move in step with the same programing language system.

We would certainly be interested in the experience of other $\mathrm{K}-12$ school systems which are using Pascal in their computer education programs. Andy, you speak of jobs for Pascal people. We are a large school systen with over 120,000 students. We invision a real demand for Pascal instructors (teachers) in the next few years, but we need help in getting started.

## Our Geraman Adress

Thank you,
IIARMSTAMI CAREER CENTEF
ATTN: SAM CALUIN
ESCHOLLERUCDKERSTR: ENDE
$8 / 6 / 5$
Samuet 10 . atru-
SAMUEL W. CALVIN
Coondinator Computer Education

Dear Aniclyğ
It was nice talking with you after having been away from the University or Minnesota for so long. As per your request, I am documenting in writing the discussion that ve had, in the hope that you will be able to communicate my request to your readers.
We are currently designing a process control language for use in our Building Automation Systems. The language will be similar to (possibly a subset of) PASCAL. In the course of our system design we have developed program in this process control language given an internal polish representation. We would like to know if any literature has been produced on the subject of decompilation rom an internal version (such as Polish or PASCAL P-Code) o a block sur topic, we would be interested in corresponding with them

Thank you very much for your assistance in helping me with this problem. Also, I would like to congratulate you on the excellent job that you have been doing with the newsletter.

Sincerely,
Dave R $\qquad$
Dave Rasmussen
DR:ph
P.S. Enclosed is my application for membership for the next academic year along with the membership fee.

April 24, 1978

Dear Andy:
I fear that in my hurry to meet the March 20 deadline for my letter I allowed two errors to slip by. They are unimportant, ut here are the corrections if anyone wants them.

Paragraph 4 line 5: delete "of".
Paragraph 4 calculations: should be
1000 transactions/day x 30 separators/transaction
$=300000$ keystrokes/day
$\cong 4 \mathrm{key}$ entry stations \& operators
$\$ 4,000 /$ month

I believe the argument is still valid
C EefndalCes
C. EDWARD REID

Kirkman Data Center
CER:jem

# : UNIVERSITY OF MINNESOTA University Computer Center TWIN CITIES <br> <br> 227 Experimental Engineering Building <br> <br> 227 Experimental Engineering Building <br> 8/12/01 

Open Letter to all PUG members from Andy Mickel

The Party is Over
I'm tired and want to quit coordinating PUG and editing Pascal News (effective any time after 1979 July 1). (* Besides, I turn 30 (base 10) on May 4 and you'll no longer be able to trust me! *)

As I said in this issue's Editor's Contribution, my ability to "coordinate" PUG and edit Pascal News may be the best I can do but doesn't seem to be enough.
is the one I mentioned in PUGiN \#12: PUG is getting too big for me to handle.

## Open Forum for Members

I can't continue; I've done all I can, and my endurance, optimism, good humor, lifestyle, physical and mental health are all stretched past the limit. The people closest to me remind me every day.

I'm also a little upset at the seemingly unnecessary growth of politics about standards, extensions, and the future of PUG and Pascal itself. The politics from my point of view seems simply a waste of time. If you reply it's inevitable, I would answer that it would have happened
conventional and ordinary manner.

## What About PUG?

What should happen to PUG and Pascal News? I don't know exactly, but there are several possibilities:

1. Disband the organization
2. Affiliate with a professional society
3. Institutionalize put
4. Keep PUG the same but decentralize the work.

Most PUG members I've talked to would like PUG and Pascal News to operate informally factually, clearly, and in a friendly manner as it has in the past. One person pointed political social or economic gain for personal benefit
litical, social, or economic gain for personal benefit.
As a fifth alternative, I somewhat doubt that anyone of you will be crazy enough to step forward and volunteer to take on all the responsibility like I did from George Richmond three years ago. Therefore I didn't include it in the list above

Disband the Organization
I occasionally have entertained the idea that perhaps the proper ending for an unconventional organization like PUG would be to simply shut it down. ("For one brie there was Camelot.") I hinted at this in the Editor's Contribution in an 12 .

Shutting PUG down might not be such a bad idea if you realize that entities can outlive their usefulness. In the long term such a decision could be considered brilliant You must realize that we have largely accomplished two important goals:

1. Making Pascal programming a respectable activity
2. Getting an officially-accepted Pascal standard.

However, my friend Steve Legenhausen told me that when the Whole Earth Catalog project was stopped, 14 -odd cheap imitations appeared, and wouldn't this also happen to Pascal News?

Actually, one of the reasons I've reprinted the roster is to provide insurance against PUG collapsing--any member had all the information necessary to restart the organization.

## Affiliate With a Professional Society

Please realize that with close to 3000 members in 41 countries, PUG functions as one of the 10 or 15 largest computing organizations in the world. We're certainly one of the most international.

That is why we are a very desirable "plum" to be annexed by the Computer Society o the Institute of Electrical and Electronic Engineers (IEEE) or by the Association fo

They have made overtures to us recently, and I asked that they put their offers in writing for publication in a future issue of Pascal News.

Although I think we in PUG do far more to promote good programming ideas and practice than do those organizations (and indeed, in spite of them!) they would offer us subscription and publication services and a guarantee of continuity. The closest example is STAPL (Special Technical Committee on APL) within SIGPLAN. But simply looking at STAPL and its

The rates for membership would surely go up, and the "membership services" would not be much better, and the publication and the group would no longer be independent. In fact you would have to join the parent organization (at absurdly high rates) or else pay nearly
$\$ 10$ more than you now pay for a PUG-only membership.

The simple fact that Pascal News would lose editorial freedom (manifested for example by our printing paper
is a major concession.

Personally, I've always been against this idea, because I never liked the way I was (and am) treated as a member and subscriber to SIGPLAN.

Unfortunately, by being late on issues such as this one, I'm not doing much better! Institutionalize Ourselves and Remain Independent

The most familiar refrain recently among PUG members besides "Keep up the good work!" and "Hang in there!" is "Keep PUG independent!".

Institutionalization would solve one big headache for me: I was never formally selected to manage PUG, because I've always considered my work "volunteer." But some people have demanded that I "represent PUG" at special conferences and make "official
statements." I've always hesitated, because at best I realized my authority was by statements." I've always hesitated, because at best I realized my authority was by default, not democratically chosen.

These people always got angry at my hesitation. (Another example of politics that's making me depressed.) I would always point to the simple mechanism of using the Open Forum section in Pascal News to air their ideas.

PUG is in good shape financially (a specific report updating the last report in PUGN $9 / 10$ will appear in PUGN 15). Recently, we have been using the extra $\$ 2$ of the new $\$ 6$ membership rate here at PUG(USA) to hire some clerical help (about $\frac{1}{4}$ of a $\frac{1}{2}$-time secretary).

But, you must recognize that technically Pascal User's Group is a non-profit activity Minnesota publication. I have taken steps all along to ensure that PUG could be transferred somewhere else within one week (really!).

The tremendous benefit we've derived from our warm University of Minnesota home should not be taken for granted. Besides paying my salary, the $U$ of $M$ has provided good production, publication, and mailing facilities. These are the major reasons the cost of PUG and Pascal News in my opinion remains reasonable. And we've, done what we have without prostituting ourselves by selling advertizing, without selling the mailing list, and without accepting subsidies from special interests (such as computer companies).

Institutionalization requires a constitution and bylaws, officers, elections, and more of the same old thing (SOT). PUG member Richard Cliche constitution and sent the following note to me on 78/08/30:
"I hope this is a good enough start. Please work over the
bylaws dues section to reflect the international situation."

- Rich (Gone Fishing!)

The proposed constitution follows this letter
In my opinion, this is the best alternative if you want to see PUG continue. However I do not want to serve as any of the officers or as the editor under a constitution and years!). The constitution would enable PUG to use authority in standards discussions and in organizing projects promoting rational programming methods. The constitution would also give, us the independence we would need to sell advertizing etc., in order to keep the cost of membership low. I don't want to waste my time making money for PUG. Count me out. I'll be the first person to step aside and not be an obstacle to the greater interest.

## Keep PUG the Same, but Decentralize the Work

This is not an alternative as far as I'm concerned. It seems that decentralization incurs the horrible tradeoff of high overhead and communications problems. If you say "nonsense" then you may be right, but then I'm the wrong person to coordinate activities scheme would be less direct work on PUG and Pascal News and more a role of an administrator.

I'm not an administrator or editor; I'm a systems programmer!
People have asked me if there was any chance that I want to do PUG full-time. The answer is "no."


I (with the generous help from many dedicated people) have had fun organizing PUG and putting together Pascal News. We've done so very informally.

I assumed the editorship after issue \#4 when George Richmond (who had edited for 2 years) gave up because of lack of time (his management was not as far-sighted regarding ADM 75 conference in Minneapolis. I was there and was "volunteered" by everyone to coordinate. So if George is Phase I, then I am Phase II.

I became involved with PUG because I wanted to see Pascal succeed, and I knew that something had to be done urgently to make that happen.

I have had fun in trying to produce a creative and refreshingly different and unconventional publication to promote a programming language. At times, it has been point back to the ordinary--the same old thing (SOT).

I assure you that the SOT approach to PUG and Pascal News would not have succeeded as well. The conventional wisdom would have doomed Pascal to the role of "just another tactics would have failed. The combination was irresistable.

Sometimes I've done things differently just to be different: such as printing paper clips and a screwdriver on the cover of Pascal News. But other differences I subscribed to. Examples: page numbers on the left in the table of contents; a single, self-explanatory POLICY; enough room on the All-Purpose Coupon to write a $4-5$ line address and comments about anything; easy-to-obtain and publicized backissues; and "all the news that fits, we print."

As for price, I would never want to be a PUG member myself if the cost of membership went over $\$ 10 /$ year '(in 1977 dollars). By keeping things simple and excluding special rates, services, etc., we have also kept the price lo
benefit of students (who show the way to the future).

Well, if you are confused, so am I! It has been sheer agony to write this letter, not because I don't want to quit, but because the ideas needed to be stated carefully.

In late October I wrote to ten or so active Pascalers for advice, and I'm grateful to Jim Miner, Rich Stevens, Rick, Shaw, Tony Addyman, Bob Johnson, Rich Cichelli, and Jeff Tobias for the advice they offered.

It has been really disappointing to be without the advice of Judy Mullins Bishop, David Barron, and Arthur Sale. They are three persons I would consider to be among the easier for me if they hal News since its beginning. I just know it would have been

But, then, that's the problem! Something else must happen. I think it's time for Phase III.

Sincerely,

The following are submitted as＇a proposed Constitution and initial set of Bylaws for the Pascal Users Group．The Cbinstitution and Byiaws will be accepted or rejected by a simple majority of the ballots（enclosed with this copy of Pascal News）returned to Rick
Shaw before April 15，1979．

A few notes about some of the wording in the documents．First，concerning the choice of an＂official＂version of the organization＇s name－－apostrophes are bad news in organizational names．The American Newspaper Publishers Association dropped one from
their name－－let＇s drop it from PUG．Secondly the term＂Chair＂is intended to be their name－－let＇s drop it from PUG．Secondly the term＂Chair＂is intended to be
equivalent to the term＂Chairperson＂．It＇s just shorter and sounds a little less clumsy．
－Richard J．Ciéhelli，August 1978

## PASCAL USERS GROUP

Official Ballot－October 1978
I believe that the PASCAL USERS GROUP：
＿＿＿should institutionalize itself and remain independent．
$\qquad$
$\qquad$ ／do Not $\qquad$ approve the submitted Pascal Users Group Constitution；
$\qquad$
$\qquad$ $l$ do Nor $\qquad$ approve the submitted preliminary Pascal Users Group Bylaws．
My reasons for rejection of either document are：
＿＿＿should NOT institutionalize itsexf，but instead should
$\qquad$ disband，or
$\qquad$

$$
\begin{aligned}
& \text { affiliate , 放ith a professional society: } \\
& \text { AGM SIG }
\end{aligned}
$$

$\qquad$
＿＿＿other：

Return this completed bailot by April 15，1979，to：


$$
\begin{aligned}
& \text { USA }
\end{aligned}
$$

Your signature need only be on the envelope enclosing the ballot．Riok will certify that voting will bee by：members ronyy：

## Constitution of the Pascal Users Group

## Article I Name of the organization

The name of this organization shall be the Pascal Users Group（PUG）．
Article II Purpose of the organization
A．：The primary objective of PUG is to promote the use of the programming language Pascal as well as the ideas behind Pascal．
B．Specific objectives shall be：
1．to provide channels of communication among members of the international Pascal community（through Pascal News，etc．）．

2．to coordinate the efforts of individuals in forming special interest groups within PUG concerned with standards，implementations，etc．

3．to coordinate sponsored research into implementations，uses，etc．of Pascal．

4．to facilitate distribution of Pascal software among PUG members．
Article III Membership
A．General PUG membership requírement
Any person who is interested in the objectives of the Pascal Users Group may become a member upon paying the current annual dues．
B．Voting rights
Formal voting privileges consist of the right to vote at PUG meetings and through mailed ballots on proposed amendments to the PUG Constitution， Bylaws，and standing rules，and on all motions．made to and by the Chair．All
members are entitled to vote． members are entitled to vote．
Article IV The Officers
A．The government of PUG shall be vested in the Executive Committee which shall consist of：

## The Chair

## The VIce－Ghair

The Secretary／Treasuree
－The Editor of the Pascal News
The most recent prevíous chaint
＂Three members－a－1arge
B．The Executive Committee members（excepting the most recent previous Chair） shatily be ellectéd for a term of two years by members of pug
c．Ant menter of RUG shatife betighte for anytoffice．The offiae of ohatim may not Te feld for more than two consecttive terms by the same individual．

Va'canciest of Office
If any office (excepting members-at-large) shall become vacant, the chai shall at the earliest possible date thereater order a special election fo the purpose of filling such office. The member thets elected shall tak office immediately and shall hold office until the next regular election.
E. Duties of the officers

1. The Chair shall
a. preside at all PUG meetings
b. call special meetings at her or his discretion subject to the
limitations of Article V, Section $E$.
d. make provision for the discharge pro tempore of necessary duties of absent members
e. sign all warrants on the treasury of PUG
g. c'arry' out assignments and instructions dictated by vote of the
h. perform other duties as customarily pertain to the office of Chair
2. The Vice-Chair shall be an aid to the Chair and in case of absence of the Chair shalil pro tempore assume and perform the duties of the office of
3. The Secretary/Treasurer shall
a. keep a record of all meetings
b. issue timely notices of meetings and agenda after consultation with
c. conduct correspondence of PUG
d. collect all fees and dues
e. maintain a list of current (paid-up) members
render an account at least yearly, or more of ten if required, of all expenditures
g. pay the bills of PUG only after approval by vote of the Executive

Members-at-large and the previous Chair shall attend Executive Committee meetings and vote on issues raised there.
5. The Editor of the Pascal News shall coordinate the publication and distribution of the journal, edit articles, and write editorials.

Article V Meetings
A. Time and Place

At least one regular general membership meeting shall be held each year, the place and time to be determined by the Executive Committee.
B. Voting

A simple majority shall be required to pass all motions. Members present shall constitute a quorum.
c. Meeting procedure

The procedure at all meetings of PUG shall be governed by this Constitutio and its Bylaws and by Robert's Rules of Örder
D. Mótions

Any member may make a motion to the Chair. This motion must be accompanied by at least one second to the motion by another member.
E. Speciall meetingis

Special meetings may be called when the Chair, after consulting with other Executive Committee members, is convinced that the need is sufficiently urgent. A special meeting shall be called Executive Committee members regardless of the wish of the Chair.

## Article VI Amendment

A. This Gonstitution may be amended at any regular business meeting of PUG by a $2 / 3$ vote of those present and voting, provided that written or printed notice amendment has and tiven to members in tictent tis for it to have been received by them at least one month before the meeting.
B. Bylaws of PUG may be adopted or modified at any regular meeting by majority vote provided that notice has been given as described above.

## Bylaws of the Pascal Users Group

article I Fees and dues
A. The annual dues shall be:

L4.00 (U.K.) per year when joining from Europe, Western Asia, or Northern
$\$ A 8.00$ (AUS) when joining from Australia or Eastern Asia;
$\$ 6.00$ (U.S.) when joining from elsewhere.
These dues are payable in advance during July
B. Members will receive all Pascal News issues of the July-June year during s, except possibly new members joining after back issues are not available.
C. Pascal News subscriptions are available to libraries and other organizations Pascal
at $\$ 25.00$ (U.S.), L15.00 (U.K.), or $\$ A 25.00$ (AUS) per year.

Article II Meetings
A. Date and time of annual meeting

The annual meeting will be held on the afternoon of the Sunday preceding the Association for Computing Machinery (ACM) annual conference at a location near the conference site.

Article III Sponsoring Affiliates
A. Individuals and organizations wishing to fund colloquia, conferences, research, and other activities of PUG may do so by becoming PUG affiliates, subject to approval by the Executive Committee

## Pipe Line Technologists, Inc.

July 17, 1978
Dear Andy,
In Pascal News \#12, J. S. Merritt wrote that he couldn't find th CACM article by Tanenbaum mentioned in PUGN \#11, p. 87. I couldn't either. As it turns out, the publication date of December 1977 is wrong. It appeared in the March 1978 issue. Here is the correct reference.

> Tanenbaum, Andrew S. Implications of Structured Programming for Machine Architecture. Comm. ACM 21 (1978), 237-246.

This is a thought-provoking article which implementors of portable Pascall systems should read. It shows the advantages of designing a computer architecture taking into account not only the formal properties of high level languages; but also impirical knowledge of how those languages will actually be used. The result is a stack machine wherein the vast majority of instructions require only one byte of code. Tanerbaum's design is called the EM-1. It could be built a a hardwired computer, microprogrammed, or--and this interests me--as a software interpreter on byte-oriented microprocessors.

The very compact object code of the EM-I will go a long way toward getting large compilers into small memories and external storage devices. Here are some code space benchmarks (complete programs) for the EM-1 contrasted with carefully handcrafted assembly language programs for the PDP-11, which is normally considered an efficient machine in code space usage:

Towers of Hanoi Sortinteger array Dot product
Find Primes PDP-11
352 bytes 992 byt
$\begin{array}{rrr}562 & \text { " } & 1,248 \\ 552 & 11 & 832\end{array}$ PDP-11/EM-1

| 52 |
| ---: | ---: |
| " | 2. 82

306 " 704 "

1. 5
2. 30

To produce an assembler and interpreter for the EM-1 machine for all the popular microprocessors would be a worthwile project. I would be happy to talk to anyone interested in the idea.
$\underset{\substack{\text { RALPH DAVIS } \\ \text { RXECUTVE DIRECTOR }}}{ }$

July 28, 1978

```
Mr. Andy Mickel, Editor Pascal News
Computer Center, 227 Exp-Engr
University of Mínnesota
Minneapolis, MN 55455
```

Dear Andy:
Pascal-ers should take note of Edsger W. Dijkstra's
rticle "DoD-I: The Summing Up" in the July 1978
article "DoD-I: The Summing Up" in the July 1978
SIGPLAN Notices, pp. 21-26. Many have been proud that
SASCAL will almost certainiy base the DoD's new
standard; the results appear likely to prove that
pride not fully justified - not because of shortcomings
in PASCAL but in the bureaucracy. To quote Dijkstra
briefly,
...instead of listing the goals to be reached; IRONMAN
aiready starts the design by prescribing "features" From which it is often hard to reconstruct or guess
which sensible goal they are supposed to serve.

And his closing,
Of ALGOL60 C.A.R. Hoare once remarked that it was a significant improvement over almost all of its successors What can we do to prevent PASCAL from sharing that fate?

Sincerely,
Cirkman Data Center
CER: jem

THE UNIVERSITY OF MISSISSIPPI SChOOL OF ENGINEERING UNIVERSITY, MISSISSIPPI 38677

Andy Mickel
University Computer Center: 227 EX
208 SE Union Street
Univ, of Minnesota
Minneapolis, MN 55455
Dear Sir,
I have been using PASCAL here at Ole Miss for the past two years on the DEC-10. I currently have available two compiler writing tools written entirely in PASCAL
(a) LEXGEN--An Automatic Lexical Analyzer Generator The generator takes regular expressions for any number lexical tokens as input and outputs the minimized inite automaton for accepting any of that set of lexical tokens. Intermediate user-controlled output includes diagrams showing how the NFA is constructed he complete NFA (in graph form), the resulting DFA (in tabular form). These intermediate outputs should especially useful for teaching the theory and pplication of this type of lexical analyzer
(b) LALR1--An LALR(1) Syntax Analyzer. Given the BNF description of a grammar this program utputs the LALR(1) tables for driving a parser. ndication is given whether grammar is SLR(1), LALR(1) r neither.
ither of these programs and their documentation is in the new ar of a specialized nature to be included I invite any interested parties to contact me directly.


AUGUST 23, 1978
PASCAL USERS' GROUP
ATTN ANDY MICKEL
UNIVERSITY COMPUTER CENTER: 227 EX
UNIVERSITY OF MINNESOTA
208 SE UNION STREET

Andy,
 old "ALL PURPOSE COUPON" and a check for $\$ 10.00$ in devalued American currency (since PUG membership fees have surely increased by now). If even $\$ 10.00$ isn't enough for two years, let me know how much more is necessary, and I'll send the balance ASAP.

I hope all is going well with pUG -- I have some doubts since $I$ haven't heard a word from you people since March. Was there another 77-78 issue published after \#11? If so, I have never received it, and I'd hate to miss anything!

A few weeks after attending the 2nd West Coast Computer Faire, I took a job as "designated internal programmer" for North Star Computers, here in Berkeley North Star is best known for its mini-floppy diskette subsystem, which is compatible with any 8080- or $z-80$ based mainframe incorporating the $S-100$ bus. To date, the firm has supported only BASIC (albeit a powerful, feature-laden version of the convinced the "powers that be" to look seriously into supporting Pascal as bot an internal software development tool as well as a marketable software product It is almost certain that Ken Bowles' group will develop a vexsion of their UCSD Pascal system which will operate on 8080 or $Z-80$ machines using North Star disk units, but it isn't clear at the moment whether or not North Star will itself support and/or market the system (though I am personally lobbying for such a development).

Regarding my somewhat sceptical comments on Pascal in the micro-world as published in PNEWS 9/10, I am pleased to note that UCSD Pascal seems to have "done the trick" and catapulted full-blown Pascal into the marketplace, at an extremely reasonable cost, yet! Finally! In the spirit of "hit 'em again, harder", part 1 of my own (after innumerable crazy circumstances and delays) in the september-October issu (after innumerable crazy circumstances and delays) in the September-October issue of Creative Computing magazine. With luck, my own, and other, similar articles will serve to bootstrap the consciousness of personal/micro-computer users into
the Pascal era. (Notice the Pascal-oriented August, 1978 issue of Byte, for example.

From my vantage point, in the midst of the small-systems market, I see Pascal's momentum increasing at an astonishing rate. It appears that we now have the ball. Let's all pull together and run with it -- now that many computerists are accepting
pascal as a real" language, there must be a concerted effort on the part of we who
for it. (Need I add that this can also be quite luc
Jim Merritt $\quad \substack{\text { for it. } \\ \text { rative:) }}$
POBox 4655
Berkeley CA 94704
Phone 415-845-4866



August 29, 1978

Mr. Andy Micke1, Editor
Pascal News
208 S.E. Union Street
University of Minnesota
Minneapolis, MN 55455
Dear Mr. MickeT :
In review of recent issues of Pascal News, I have noted several letters from readers in regard to Pascal jobs. Most expressed amazement in their success in finding PASCAL positions.

I thought it would be of interest to you and your readers that Computer Careers, Inc. Agency has a full division of consultants working with. PASCAL type programmers. The demand for the higher leve] block structure languages is growing everyday. We have been quite fessionals in their pursuance of PASCAL careers.

If we can be of any help to you or your readers, please fee]. free to call.

$\mathrm{CB} / \mathrm{r}$
Enclosures

## INFORMATION ENGINEERING COURSE

DIVISION OF ENGINEERING
UNIVERSITY OF TOKYO GRADUATE SCHOOL

## Sunkyoku, Tokyo 113 Jan

Pascal Users Group
c/o Professor Arthur Sale
Department of Information Science
University of Tasmania
BOZ 252C GPO, Hobart, Tasmania 7001
Australia
Dear Professor Sale:
Enclosed please find our renewal remittance \$A56, for the Pascal

Users Group membership 1978-1979 for seven of

## M. Arisawa

T. Hilkita
S. Yoshimura
N. Tokura
M. Takeichi and

Our addresses remain unchanged. As to other Japanese members, Messrs. H. Ishida, M. Watanabe, K. Noshita, N. Wakabayashi and H. Nishioka have is presentreaty or paid more than one year's members fee. Mr. Kishimoto

I am so sorry for not writing you earlier. We are one of the first Group who introduced pascal in teaching programming. In my class, all the examples were swithed to Pascal since the fall semester of 1972 , and the first Pascal compiler beame available in the summer of 1974 .
Since then at the University of Tokyo, three versions of pascal compilers Since then at the University of Tokyo, three versions of Pascal compiler haboratory, a pretty printer for Pascal has just been completed. The pretty printed output is obtaind through the phototypesetter which really generates very high quality documents. Besides this ${ }_{r}$ we are stili. considering of rewriting the Pascal, report. in more accurate and understandable way. The Pascal compiler in Pascal may be improved to become much more Pascal lijke, that is, with fuller, Pascal spirits.

I hope we are able to see each other at the IFIP congress two years later, in 1980 .


Einti Wada Professor
years later, in 1980 . Sincerely yours,

Users Gr
$\qquad$
$\because$
EW/mk
enc.

Health Products Research, Inc.
3520 U.S.Route 22, Somerville, New Jersey 08876 • (201) 534-4148
Cable Address: "Healthipro"
23 September 1978
Mr. Andy Vickel, Editor
Pascal Hews
University Computer Center: 227 EX
208 SE Union Street
University of Minnasota
Lear Andy:
Às 'Pascal Coordinator' for the Amateur Computer Group of New Jersey (ACG-NJ), I an in a position to report some good news about the enthusiasm for Pascal among computer hobbiests in the New Jersey area:
a. The ACG-NJ has taken advantage of the group
subscription of $\hat{\text { fer }}$ of Ken Eowles' group at the University of Subscription offer of ken Lowles group at the miversity obtained the UCSD implementation throush the ACG, and at least five have it "up and running" on their personal systems. Most of these systers are $3080 / \mathrm{Z}-80$ microcomputers, although there are two or three LSI-1!s as well
b. I gave a brief talk on Pascal at this month's aCG-HJ meeting, which was well received; I have also been invited to speak on Pascal to the New York City amateur conputer group in jecember.
c. ft least sixty people attended a "Pascal lisers Group" session at the "Personal Computing "73" show; helc in the Philadelphia Civic Center at the end of Aupust. This turnout was mildily astonishing in view of the fact that the session on Pascal vas a last-minute addition to the program, not show, and scheduled on the show's last day. A show of hands at the start of the session yielded the following statistics:
Persons who had used UCSD Pascal:

Persons who had used another Pascal:
16

Persons who wantca but did not have Pascal: 25
Persons who didn't know whether they
wanted Pascal or not:
Persons who knew they didn't want Pascal: Non
I wish to thank lir. Fobert Hofkin of UCSD, who happened to be at the show on business, for stopping by and helping field some of the questions.
d. Three other noteworthy presences at "PC 78 " were those of three companies selling UCSD Pascal with their computer systems:
(1) Northwest Microcomputer Systerns, of Eugene, Orgeon, demonstrated their "Programmer's Workbench", a desk-top system containing an 8085 microprocessor running at 3 (optionally 5) $\mathrm{MHz}_{,} 56 \mathrm{~K}$-bit bytes of memory, dual magnetic diskette drives and a video display, priced at $\$ 7495$. I understand that two people from Zurich wanted to pay cash and walk away with one wouldn't have run on Swiss electrical power without a modification too extensive to be done at the show), and that Carl Helmers, editor of BYTE magazine, was responsible for the disappearance of one of the systems on the second day of the show, being "unable to resist the desire to take it to his hotel room and play with it".
(2) ALTOS Computer Systems, of Santa Clara, California, demonstrated their "ACS8000" system, featuring a $\mathrm{Z}-80$ microprocessor running at 4 MHz , up to 64 K by.tes of memory, system with 32 K bytes of memory and a single drive: $\$ 3,840$. (This system does not include a built-in videa display.)
(3) Alpha Microsystems demonstrated UCSD Pascal as a subsystem of their multi-user system, whose CPU is based on the Wester Digital WD-16 chipset. Workspace available to a bytes, but the system supports multi-megabyte hard-surface magnetic discs as well as (or instead of) diskettes. I regret I do not have their prices readily at hand.
e. The August issue of BYTE magazine had a cover
portraying "Pascal's Triangle", an area of smooth water with well-marked channels bordered by such less hospitable places as the turbulent "FORTRAN Ocean", the desolate "Isle of BAL" much commercial traffic is seen), the perpetual fog bank wherein lie the "exotic and mysterious jungles of LISP", and the "interactive and weed-filled Sea of BASIC". Several and vessels, ranging from warships to tiny rafts, are fleeing to the safety of the Triangle.

On a more serious level, the same issue of BYTE containe five articles on Pascal, including one by Ken Bowles himself entitled "PASCAL VERSUS COBOL: Where Pascal Gets Down to since there seems to be a consensus among those involved in the "personal computer" industry that the big market right now is small business systems, for which the greatest lack is high quality software.
A less welcome development is the discovery that the UCSD is no more inmune than any other vendor to the announce-it-early, for their Release I. 5 , my phone calls every other week being taken by a pleasant but apparently not-too-knowledgeable you person who assures me that the Release will be forthcoming "in another week or two". I guess we should be thankful we get anything at all.
keep up the good work!

P.S.: I prepared this letter on my personal systern using the screen editor that comes with UCSD Pascal. It works!

July 10, 1979

Andy Mickel
Pascal Users Group
Univ. of Minnesota
Dear Mr. Mickel:
Enclosed is my renewal for the coming year. I have trully enjoyed receiving PUG newsletter. (I finished \#12 in less than 12 hours and still haven't read GACM May!) The new section on APPLICAATIONS should be an excellent media for transmittal and evaluation of programming methods.

We at Fitchburg state College have just totally restructured our counse structure to put Pascal into the Freshman year where it belongs. Other languages are taught within the courses which require them and assume a Knowiledge of Pascal.

Re: standardization of Pascal. I vote for Charles Fischer's method (PUG \#12, pg. 54), a standardized set of extensions designed by a small group and an all-or-nothing vote by PUG membership. I have a great many changes I'd like to see in Pascal; but, I'frather see a standard. I'm sure
a lot of other people feel the same way.

Keep up the excellent work.


Computer Science Program
Fitchburg State College

MEAR ANIY,
THERE TS A NEETIFOR A BOOK TO BE FUBLTBHELABOUT FASCAL, TT WOULTH HAVE TO WEFTE A GTANTABL SUCH AS WTETHS (7)" HAVE A COMFLETE LTSTTNG OF THE COMFTLEF
 WTH HABTMANNS (3) COMES CLOSE TO THE THEAL ITT WCULT RE ON THE OFDEF OF

 ON A MYEOY MNT, OE MAXT-COMFUTER FABCAL PQ TN WTRTHY ( 7 ) NN NOT COMFETE, (MTESTE GENERATEG COLE AND WRTTEN TN A HTGHER LEUEL THAN FO). FASCAL-S TN

WTFTHys(\%) ALSO FAS MTSSTNG GENEFATEN CONE FASCAL FA FERHAFS IS COMFLETEy BUT BUT THESE THEEE LEVELS AFE AFPROXIMATELY WHAT IS NEEMED,
1.) HALSTEAT: MACHINE TNDEFENDENT COMFUTER FROGRAMMING
2) HANEEN: ARCHTECTURE OF CONCUREENT FROGRAMS
3) HARTMANNS: A CONCURRENT PASCAL COMFTLER FOR MINI-COMFUTERS
4) MCKEEMAN: A COMFTLER GENERATOR
6) WATTE" MMPLEMENTNG SOFTWAREMFEMENTATION
7) WATTE A MFLEMENTNG SOFTWAEE FOR NON-NUMERTC AFFLICATTONS
B) WTETH A NENEN F'ASCAL USER MANUAL \& REFOFT
8) WIRTH: ALGORTTHM + MATA STRUCTURE = FRGGRAM

PASCAL User's Group
C/0 Andy Micke1

## Dear: Andy:

At Computer Automation's NAKED MINI Division, PASCAL is gaining
interest and support. Our compiler on DOS4 produces code for a interest and support. our compiler on DOSA produces code for a
virtuall machine. I have recently converted the machine to work under our new operating system OS4 on the NM4 series computer. The same compiler now runs under DOS2, DOS4 and OS4. For
marketing information, contact Laura Cvetovich (M/S 1167).
PUG members might be interested to learn that several openings in system software development are available at CA requiring PASCAL and assembly language experience. The inside track
can be had by writing Dave Robertson (M/S 1175).

Keep up the good work,
D. J. Maine

Research Scientist
M/S 1175
: . . . i
informatik
Prof. Dr. H.-H. Nagel

## CALIFORNIA INSTITUTE OF TECHNOLOGY

## Division or mioLoay 210.7e

Andy Mickel
ditor, Pascal News
University Computer Center
227 Exp. Eng. Bldg.
finneapolis, Minnesota 55455

Dear Andy,
Judy Bishop's discussion of subranges and conditional loops (Pascal News \#12, pp 37, 38 and 51) clearly states a basic problem in standard Pascal: how to keep index variables within their subrange at all times. However, her solution does fixes would hinder readibility; (2) the necessary extra type definitions are a hassle both to write and to read; and (3) the extra allowed value of the index may in some cases degrade the ability of the run-time checks to stop an error at it's source.

Some Pascal compilers, such as the Brinch-Hansen DEC-10 compiler in use here at Caltech, allow the loop. . exit if construction. Loop... exit if is the most general form of the conditional loop, since it contains a statement block before lock after it (as does while). This generality is necessary for a natural sol tion to some problems, including this one.

Having defined $i$ on the subrange min..max, we can write:

$$
\begin{aligned}
& i:=\min ; \\
& \text { loop (* something *) }
\end{aligned}
$$

$$
\begin{aligned}
& \text { exit if }(i=\max ) \text { or (condition); } \\
& \underbrace{}_{i}=\operatorname{succ}(i)
\end{aligned}
$$

end;
Thus we always have $i<=$ max. In standard Pascal, a somewhat less elegant but equivalent construction is available:
i : = min;
$\frac{\text { while }}{(* \text { true }}$ some $\frac{\text { to }}{\text { thing }}$ *)
if ( $i=\max$ ) or (condition) then goto 10 ;

$$
i:=\operatorname{succ}(i)
$$

end; (* next statement *)
yours sincerely,
Kail Fryuell
Karl Fryxell

## Dear Andy:

I'm enclosing twelve (148) dollars for two years PUG dues. If you don' think PASCAL will survive for that long please return some of my dues.

Some general comments. First in response to "What To Do After A While" you need two new operators
a AND THEN b; C OR ELSE $\quad d$. In a more serious vein, many PUG articles contain phrases like "stamp out
FORTRAN", or "kill the dinosaur", etc. The articles seem to be written with all of the grace and charm of a stiff necked missionary trying to convert a bunch of ignorant heathens. Why is there such an emotional investment in promoting PASCAL? PASCAL, like most human inventions, has some good points and some bad points. PASCAL was implemented on a CRAY-1 Computer by a group at Los Alamos. There also exists a group of 18 short "kernels", called the Livermore Kernels, which (allegedly) are typical of the kernels ran (last January) with an average "speed" of about 3.6 MFLOPS (million floating point operations per second). If a second program is (million floating point operations per second). If a second program is When run using the current CRAY FORTRAN Compiler the rate is about 22 MFLOPS, planned FORTRAN enhancements (for "this year") should bring it to ver 30 MFLOPS. Now, there are significant differences in implementation strategies between the LASL PASCAL and CRAY FORTRAN and it would be very wrong to conclude (from this example at least) that PASCAL is not a good language. However, with performance ratios of between 4 and 10 (depending
on one's point of view) on a sysitem that costs up to $\$ 9$ million, it seems just as wrong to conclude that "FORTRAN is obsolete". If PASCAL is to ecome a universally used language won't implementation become machin dependent with additions (and deletions) to take advantage of particular hardware?

Page \#2

This leads to the second point. I understand that there recently was a first (annual?) PASCAL standards meeting. I've heard from two different obnoxious) that nothing in PASCAL should be changed, wirth has spoken and not a "," must ever be changed. Is this realistic? As the language is used shouldn't it grow, much like English or FORTRAN when deficiencies are discovered? AFTERALLTHEREAREFEATURESINPASCALWHICHDONOTNECESSARILYMAKEA PROGRAMEASILYUNDERSTANDABLE.

More important, is this a legitimate stand for a "PUG representative" to take. I'm a PUG member and I don't recall ever being asked whether or no PASCAL should be changed. Certainly everyone is entitled to an opinion
about the future of PASCAL but shouldn't a "PUG representative" somehow survey his members? It seems to me that most of the articles in PASCAL NEWS deals with proposed additions or deletions and most of the implementations mention deletions.

I hope this gives you something to write an editorial about
Best


Richard A. Hendrickso

## RAH:al

(* In a phone call to Dick in early October, I (Andy) thanked him for the letter and explained that one reason that Pascal is at a disadvantage when compared to FORTRAN s because of the vast difference in the person-years put into compilers, libraries tc. However, I appreciated the data he provided and his feelings. I also told extensions, and that the so-called obnoxious PUG representative was Richard Cichell who indeed upset many persons. Since I wasn't there, I can only repeat the reports I have heard. I explained that human languages and programming languages are vastly different, and no, programming languages shouldn't necessarily grow, and in fact Edwin Newman's recent books (one is Strictly Speaking) deplore the unnecessary growth in the English language. And if you have a decent Pascal implementation

## Department of Applied Mathematics and Computing Science

Professors
D N de G Allen, W D Collins, S C Hunter, J R Ullmann

Andy Mickel,
Editor, Pascal News,
4th September, 1978
University Computer
208 SE Union Street,
University of Minnesota,
Minneapolis', MN 55455, U.S.A.
Dear Andy,
My copy of Pascal News $\neq 12$, mailed in Minneapolis on June 23rd arrived on August 31 st . In it I read that the publication deadline date for $\not \neq 13 / 14$ was August 15 th, 16 days earlier! Can this be true? I hope not. Please try to squeeze in the enclosed paper "Know the state you are in". Although written in great haste it might solve a few problems for a few people or at least shed a new light on them.

I was interested to read Judy Bishop's comment about booleans (Pascal News $\# 12$, page 51). Since first teaching Pascal three years preference to booleans and if. The programming style in my enclosed paper is a natural consequence of this. I have also been following recent work by experimental psychologists studying the (detrimental) effects of negation in programming logic and, in particular, the negation implicit in else. One consequence was that I submitted a paper to CACM in March of this year supporting two-state scalars and case in preference to booleans and if and, of course, praising Pascal for encouraging this approach. So please note, Judy, your anticipated
paper "Booleans considered harmful" has already been written!

$$
\begin{aligned}
& \text { Yours sincerely, } \\
& \text { Laurence OXtepinse. }
\end{aligned}
$$

L.V. Atkinson

## University of the Witwatersrand, Johanmesburg

DEPARTMENT OF APPLIED MATHEMATICS
1 Jan Smuts Avenue, Jonarnesturg, 2001, South Africa
Telephone 39-4011, Telegrams 'University', Telex 8-7330 SA

Mr. T.M.N. Irish
5 Norse Way,
Sudbury,
Chepstow,
Gwent NP6 7BB,
United Kingdom
tolephone ext
our reference
our reference JMB/SW
date 27 September 19.78

Dear Mr. Irish,
Many thanks for your note on "What to Do After a While". I would just like to clarify our points of agreement and disagreement and just like to clarily our points of agreement and disagreement and boolean operator contraversy can now die a natural death.

1. We take your point about potentially undefined factors. What ycu are saying is that the a $[i]$ in (a) below is permissable because it is an expression, but the a [i] in (b) is not, because it is a factor.

$$
\text { (a) if } i<=n \text { then if } a[i]=\ldots
$$

$$
\begin{aligned}
& \text { (b) if }(i<=n) \text { and (ali] }=\ldots \text { ) } \\
& \text { this may be a valid distinction }
\end{aligned}
$$

stinction, it is a hard one to grasp After all, the a[i] in (a) starts off as a factor!
2. TO POP-2 and RTL/2 remember to add Wirth's new language Modula and Euclid. All of these specify that factors in a boolean expression ill only bay be old, but it is certainly no longer strong in the world of language design.
3. The andop function is "wrong". Moreover, the loops in our Appendix The andop function is "wrong". Moreover, the loops in our

We could argue on and on about this for ever. Fortunately, the problem - that of searching a list to our satisfaction - has been solved in a completely novel way by Jaurence Atkinson of Sheffield University. He brought to our notice the following solution which takes account of the fact that there are three states in the loop, represented by

$$
\begin{aligned}
i<=n \text { and } a[i]<s i t e m & : \text { scanning } \\
i<=n \text { and } a[i]=\text { item } & : \text { found } \\
i>n & \\
& : \text { notthere }
\end{aligned}
$$

Solution 4. USE A STATE VARIABLE
var table : array [l..maxsize] of whatever; state (scanning, found, nothere)
$\vdots$
index:=I; state:= scanning;
repeat
if index $>$ maxsize then state:= nothere else
if table[index]<= item then state:=found else index:=index+1

$$
\text { until state }\rangle \text { scanning. }
$$

It may not be as short and sweet as your favourite solution but it works for all cases and does not need additional elements. Incidentally, this method still requires index to be declaxed over l..maxsizeplusone. (See Mullins PNi2 (1978)"Subranges and Conditional Loops").

I think we should let this matter rest now. In a sense no-one has won - we can't have undefined factors, you can't have side effects in functions. Pascal is a double edged sword, but it is very sharp

Best wishes,


Judy Bishop

Please direct all enquiries for this section to Tony Addyman Dept. of Comp. Sci.
Univ. of Manchester
Oxford Road
Manchester, England
Manchester, England
M13 9PL U. K.
or Rick Shaw Systems Engr. Labs
6901 W. Sunrise Ft. Lauderdale, FL
33313 USA

33313
Much has happened since issue \#12 last June. Rick Shaw is now Tony's "right hand" in the USA. Thanks to Tony and Rick, Standards discussions are placed within the Pascal User's Group where it belongs. Arthur Sale was selected to chair the Australian Standards Assoc. committee on Pascal Standards (MS/20).

On 78/06/18, Niklaus Wirth wrote that there was one error in the EBNF syntax published in Pascal News \#12, June, 1978 on page 52. The definition of FieldList should be:

$$
\text { FieldList }=\text { FixedPart [";" VariantPart] | VariantPart }
$$

Below are reports from Tony, Rick, and Brian Wichmann. Rich Cichelli reported that when they are ready, he will distribute the Wichmann-Sale Validation Suite and a standardsconforming checking program. (There exists a similar program developed for Pascal programs accepted by the Pascal-P compiler.) Rich expects that he will be able to distribute this software for a reasonable fee.

Tony's working group produced a third draft of the BSI/ISO standards document which will appear as PUGN \#14 (January, 1979). The BSI/ISO standards effort, incidentally was unanimously endorsed by the participants of the UCSD workshop on extensions in July. On October 11, it was reported that the ISO vote on the BSI proposal was 8 in favor (U.K., (with qualification); 1 opposed (Japan), and 10 abstentions (!?). Also in October ANSI (wnounced the formation of X3Jg, a committee for examining the ISO standard to be adopted as an American standard.
Rick Shaw, Rich Cichelli, and Jim Miner will attend as PUG's official representatives to the December 19 meeting.

## News from the International Working Group on Pascal Extensions

In PUGN \#12 we announced the formation of this group: a small number of competent
implementors of "major" Pascal implementations were chosen by Niklaus Wirth. Why only
with everyone interacting it won't work. We must rest assured that if someone has a great
idea, it will certainly be recognized. Even with a few people, it has been an overwhelming amount of work (forests of paper have been consumed!). So the project is delegated to a small group for good or ill. Here is the invitational letter from Niklaus:

एax EIDGENOXSSISCHE TECHNISCHE HOCHSCHULE
zÜRICH
Institut für Informatik
January 30, 1978

Dear Andy,
The "Standardization of Pascal" is a recurrent theme. As you probably know, I have been rather reluctant to get involved in such an effort, being aware of the time-consuming nature of ill I am also av "standard".
recent visit of Professor Jorgen Steensgaard-Madsen from Copenhagen, implementor of Pascal for the Univac 1100, has brought up the topic again. We have had some refreshingly productive discussions. The gist of them is that we should try to obtain a consensus among a few implementors of pascal Their agreement to work on such a consensus and to implement the results on their machine would in our opinion be the most effective way to reach a standard that does not only exist on paper and evokes a lot of discussion and controversy, but will effectively be adhered to:,

Jorgen has agreed to work out a draft of a working document within the next two or three months. We are solliciting your suggestions. If a positive response should emerge, we would envisage a meeting, preferrably sometime this summer. I would appreciate to know your reaction to such a plan.
The draft document to be worked out rests on the basic assumption that Pascal as defined by the Revised Report shall essentially remain unchanged. It shall concentrate on three topics:

1. Standard representation of programs in terms of standard character sets, and definition of the set of standard procedures, types, etc., 2. Clarification of issues that are left open by the Report (such as type equality), and 3. Extensions. We agreed that the following topics would be included:
2. Specifications of the types of parameters of formal procedures. This would be the only point involving an actual change of pascal, since it would require that such types be specified.
3. Array parameters, especially the possibility of omitting the Array parameters, especially the possmal arrays. This might or might not include dynamic actual arrays.
4. An "otherwise" clause in the case statement
5. Structured constant definitions.
6. External procedures and "forward" declarations.
7. Standard procedures for reading text files according to the program schemata used for regular files.

I am looking forward to your reply and suggesions and hope that with your dedicated help a contribution towards a much discussed goal may evolve. please send a copy of your reply directly to Jorgen.

Sincerely yours,
Neklaus
Prof. Niklaus Wirth
cc: O. Lecarme, Université de Nice, France (CII)
A. Mickel, University of Minnesota, USA (CDC)
H.H. Nagel, $\begin{aligned} & \text { J. Steensgaard-Madsen, University of } \\ & \text { Copenhagen, Denmark (Univac) }\end{aligned}$
J. Steensgaard-Madsen, University of Copenhagen, Denmark (Univac)
J. Tobias, Australian Atomic Energy Commission, N.S.W. Australid
J. Welsh, Queen's University of Belfast, North Ireland (ICL)

Jim Miner and I suggested in February in our response to this letter that Arthur Sale,
Tony Addyman, and Ken Bowles be added to the list because Arthur's Burroughs B6700 and Ken's microprocessor interpreters were major implementations, and Tony had been doing all of the standards work so far. They were added. We promised in issue \#12 to report on the results. In the 3 months of activity (from April to June) no one would have predicted the amount of controversy and heap of paper generated by the 10 participants It is an example of the "frailty of human interaction as opposed to problems caused by individual personalities." Nevertheless the Working Group rebuffed the hack changes done by individual implementors by concentrating on just a very few issues. We finally agreed
on some results. Our first result involved a conventional form for the almost universal extension providing an "otherwise clause" to the case statement. Arthur Sale presented the report below for publication

## International Working Group on Pascal Extensions

## Consensus Position on Case defaults

1. Background

The International Working Group is a group of implementors of Pascal et up by Nihlaus Wirth and the Pascal Users Group to responsibly draft some key extensions to the programming language pascal. The following report details the first consensus decision by the Group, and is published in Pascal News in the interests of other implementors and to achieve rapid dissemination of information.

The term conventionalized extension is used here to mean that the feature described is not to be considered as part of the standard language Pascal but rather that some implementations may include the feature in accordance with the conventions suggested by the Working Group. The purpose of conventionalizing extensions is to
(i) enhance portability of programs which use the extension, and
(ii) ensure a concern for the integrity of Pascal in making extensions.

The following minor extension to the language is the first consensus decision by the Working Group and is to be regarded as a conventionalized extension.
2. Notation

The modifications to the syntax will be described in EBNF notation, as this is likely to be the form used in the draft standard for Pascal, and an be used to avoid repetition or the introduction of new non-terminal symbols.
3. Purpose

The extension described allows a construction to which control is transferred if the selector expression of a case statement fails to match any case writing of lexical analysers so as to ensure robustness against unexpected input.

## 4. Modifications to the Report

(a) Add to the list of special-symbols in section 3:
| "otherwise"
(b) Replace the production for case-statement by:
case-statement $=$
"case" expression "으"
case-list-element \{ $" ; "$ case-list-element \}
[ "otherwise" statement \{ ";" statement \} ] "end"
(c) Add the following text to the explanation of the semantics of the case statement in section 9.2.2.2.
"If there is no constant in the case statement whose value is equal to the current value of the selector, then the group of statements between otherwise and end are executed. If the otherwise part does not occur, then programs which cause this to occur in execution are invalid."

## 5. Implications for variant records

he working Group considers that no corresponding change should be made in the symtax of variant records
6. Considerations taken into account

In recommending this syntax and semantics, the Working Group has (a) many alternatives, including
(b) the use of alternative word-symbols, including else,
(c) what the 'undefined' actions might be, and
(d) whether the extension was needed and added to the power of the language.

NOTE
creation of a draft standard, the wording of the Revised Report ay be altered, with consequent effects on the phrasing of this extension note. The syntax and semantics will not alter.

The full specification of all parameters to procedures and functions which are thenselves parameters was agreed on. Discussion of this topic was very influential and resulted in its inclusion in the third working draft of the BSI/ISO Pascal Standards document. Its description appears in \#14, so we won't waste room here.
In July the UCSD Workshop referred important extensions to the Working Group, such as conformant array parameter bounds

On August 24, Jorgen Steensgaard-Madsen had to resign as coordinator of the Working Group because he began spending a sabbatical year. Charles Fischer of the University of coordinating the Working Group. The current topicmentations. Jim Miner and I are now array parameters, which are important for building practical subprogram libraries for both numeric and non-numeric applications. - Andy and Jim
 Institute. As part of that effort, I am collecting together (with

Dear Andy
PASCAL test suite
12 June 1978
Readers of PASCAL News will be aware of the standardization effort that is being undertaken in the UK under the auspices of the British Standards

## hat

Arthur Sale and others) a suite of test programs designed to illustrate trouble spots in the language definition (and potentially in compilers). When the standardization is completed, it should be possible to use the suite to validate compilers, assess their performance or diagnostics as well as giving some indication as to how they match up to the standard.
Anybody who would like a copy of the tests or who would like to contribute to the tests should write to me. I am not publishing the tests at this juncture since they will change rapidly over the next year.

Yours sincerely
Brióm Wwhtuann


NATIONAL PHYSICAL LABORATORY
Teddington Middlesex TW11 0LW
Telex 262344 Telegrams Bushylab Teddington Telex
Telephone 01-977 3222 ext

This report will necessarily be brief, since time spent writing the report is time that cannot be spent on the draft
the April meeting of DPS/13/4 it was decided that we should make an attempt at preparing a draft. Up until the April meeting our efforts had been largely dirocted towards identifyinc the problems rather than the solutions. Although production of the rough draft was rather behind schedule (largely due to examination marking by the university members of DPS /13/4) I was able to take a copy to LaJolla. This proved to be very valuable. An improved draft was presented at the September meeting of DPS/13/4 at which a number of When completed, this working document will be qiven to BSI for the necessary editorial and other processing before it is issued as a draft for public comment.

When I was in the USA this summer it was my belief that I could arrange for the draft $B S$ to appear in Pascal News. This will not be possible, unfortunately. However, I will be submitting to Andy
Mickel a copy of the working document which DPS $/ 13 / 4$ passes to $B S I$. The tor unless any errors are detected and corrected by the bis machinery.

The decision to prepare a draft for public comment does not mean that we have, or even believe that we have, resolved all the questions that people have concerning Pescal. We lave-prepared a draft because we believe that many issues have been resolved and that now is an opportune time to receive comnents on what has been done.
It is my intention to send a commentary on the working document along with the working document, in an attempt to highlight those areas which currently are causing concern.

In the next issue of pascal News I should be in a position to report on the situation within ISO.

A M ADIYMAN
AM Xddywen
onvenor - DPS/13/4

## SYSTEMS

Mr . Andy Mickel, Editor
PASCAL News
208 Southeast Computer Center
208 Southeast Union Stree
Minneapolis, Minnesota 55455
Dear Andy:
I apologize for writing this letter so late. My only excuse is that I have been quite busy
As you know by now, Andy, during the conference at the University of California, San Diego, I volunteered (read I was cajoled!) to least coordinator for standards for the PASCAL User's Group at officers and the like.

However, after conversations with Tony Addyman, we both noticed a severe overlap in his informal position with the User's Group and mine. We solved this quite easily. for that. So here is what I offered to do:
(1) Act as North American liason for Tony's efforts in standardization and to generally aid him and the newsletter staff.
(2) Draw up a first-draft proposal for Program Interchange Standards (by January 1, 1978)
(3) To collect and standardize a more extensive set of syntax and semantic test programs for the standard PASCAL language these tests (by April I, 1979)

I volunteered to do the last item because it is one of the things I committed to do for my company. I know that both Arthur sale and will work closely with them as well as any other Pascal User's Group members who wish to contribute.

I hope to be able to carry out these tasks in a timely manner.

WFS/esI

W. F. Shaw

UNIVERSITY OF MINNESOTA twin cities
William F. Hanrahan, X3 Secretary American National Standards Institute, SPARC Suite 1200
Washington, DC 20036

University Computer Center
27 Experimental Engineering Building (612) 373.4360

Wednesday 78/09/27

Dear William,
1 am writing to you regarding an article appearing in the 28 August issue of Computerworld on page 27 about the proposed ISO Pascal Standard. I only hope that this letter is not too late.
I.am the coordinator of the international Pascal User's Group (PUG) which now numbers over 2700 members in 41 countries and 49 states. PUG produces the quarterly Pascal News, of which 1 am editor. Pascal News has overseen the rapid spread of Pascal simply by disseminating vast quantities of information (please see enclosures) American PUG members number about 2000.

I would like to point out that PUG has been in very close contact with the British Standards Institute DPS/13/4 group and its chairman, Tony Addyman, which are producing roup the Swedish Technical Committee on Pascal, the French AFCET Sub-group on Pascal and the Pascal group within the German ACM

The PUG membership (which certainly comprises the majority of Pascal enthusiasts) has consistently held the position that the control of the international effort be left in British hands.

The over 100 persons attending the Third Annual Computer Studies Symposium on Pascal at the University of Southampton, held in March, 1977, unanimously approved a with mox

Niklaus Wirth, the designer of Pascal, has given Tony Addyman his enthusiastic support and is providing technical assistance for the ISO standardization effort.

The recent Pascal Workshop held at the University of California, San Diego, was attended by representatives of over 15 computer companies having used Pascal for more than a year. That workshop unanimously agreed that every attempt be made to confo
to the resulting BSI/ISO standard.

We are proud to say that the simplicity of Pascal which separates it from other anguages has carried over into the standards activity undertaken so far.

In a couple of issues of Pascal News, people who supported a general standards effort naturally thought of turning to ANSI. It was pointed out, however, that the ANSI Pascal Standard should be one line which reads: See the ISO Pascal Standard document number X." just

Thus ANSI has the opportunity of reciprocating its respect with ISO. Pascal is a language with European origins, and the major work on standards has appropriately been left to Europeans. The savings in time, expense, and energy to ANSI or any proposed American "technical committee for Pascal" are obvious. We all don't want national variants for Pascal

american national standards committee
X3-computers and information processing $\times 4$--office machines and supplies

## Aprating under the procoduros of the

For more information, contact:
C. A. Kachurik

202/466-2288
PASCAL PROGRAMMING LANGUAGE STANDARDS COMMITTEE FORMED BY X3
Washington, D.C. - "Programming Language PASCAL" is the responsibility of a new committee under American National Standards Committee X3. Identified as Technical Committee X3J9,
the initial task of the technical committee is to prepare a proposal for standardization of the PASCAL programming language and obtain approval of the proposal and program of work. Committee work will be aligned closely with the international standards subcommittee on PASCAL, as well as on-going work in the Federal Government, domestic professional societies, equipment manufacturers, and other interested organizations.
The committee is seeking active participation from users of PASCAL, as well as developers of the PASCAL language compilers. Interested persons and organization representatives are invited to contact C. A. Kachurik, 202-466-2288 at CBEMA/Standards, Washington, D.C. for rther details.
X3J9 will be a part of the parent committee X3, which has overall responsibility for standards on computers and information processing. X3 currently has 29 technical committees and has completed some 60 standards published by the American National Standards Institute secretariat is the Computer and Business Equipment Manufacturers Association (CBEMA)

Poter
rican national standards committe
$\times 3$-computers and information processing
X4-office machines and supplies

## Amarican National Standards Institute

NEWS RELEASE
November 10, 1978
For more information, contact:
C. A. Kachurik

PASCAL PROGRAMMING LANGUAGE STANDARDS COMMITTEE MEETING SCHEDULED
Washington, D.C. -- Mr. Justin Walker will convene the inaugural meeting of the newly-formed X3 Technical Committee on Programming PASCAL, X3J9, Tuesday, December 19, 10:00 a.m. at the offices of the Computer and Business Equipment Manufacturers Association. The Association T.W. Suite 1200 , extensive background in the PASCAL area and has developed several compilers.
The committee is seeking active participation from users of PASCAL, as well as developers of the PASCAL language compilers. The initial task is to prepare a proposal for standardization of the language and obtain approval of the proposal and program of work.
. C. A. Kachurik, Secretariat Staff, at 202/466-2288 for further details.


Special note: We are pleased to print Scott Jameson's announcement (below) of the formation of the PUG Implementors Group. Given the wide variety of previous and current actuicies in the in leranion an read and respond to Scott's proposal.

As this is the first issue of Pascal News in this academic year, let us explain how this section is organized:
-- First, Reports of interest from the Implementors Group.
-- A CHECKLIST to be used as a guide to users, distributors, implementors and maintainers for reporting the status of Pascal implementations on various computer systems.
-- A PORTABLE PASCALs section reporting distribution information about kits used to produce Pascal compilers for real computer systems.
-- Information on PASCAL VARIANTS.
-- A FEATURE IMPLEMENTATION NOTES section describing implementation strategies and detaills of various Pascal features as suggestions to all the compiler Implementors Group's Reports.
-- A list of MACHINE DEPENDENT IMPLEMENTATIONS sorted by name of computer system, giving news of Pascal compilers for real machines.

- And an INDEX to all the implementation information in current issues and back issues of Pascal News.

Note: It is not economically feasible for us to reprint all of the old information from previous issues. We therefore will provide references to back issues when we have received no new information. (Use the All-Purpose Coupon at the beginning of this issue for ordering back issues.) We will be very happy to print new information, or revisions of previous items, submitted by users, distributors, maintainers, or implementors. When appropriate please use the CHECKLIST form. We prefer dark camera-ready copy,
single-spaced, with wide ( 18.5 cm ) lines.
IMPLEMENTORS GROUP REPORT

PACKIARD
DATA SYSTEMS • 11000 Wolfe Road, Cupertino, California 95014, Telephone 408-257--7000, TWX 910-338.0221

TO: PASCAL NEWS Readers

One of the results of the UCSD Workshop on Extensions to
Pascal was the decision that the Pascal User's Group would evolve in order to continue to meet the needs of its $2200+$ members. The
Implementation Notes
programmers to language designers, require PUG to structure itself so that it can better respond to everyone interested in to provide a medium for communicating items of interest for those involved in developing pascal compilers, and for those desiring further information regarding specific compilers or machine implementations.

The Implementors Group can serve the pascal community in many ways. Some of the things we hope to do are:
~ Publish a newsletter, aimed primarily toward the compiler developer. It will contain articles of interest to implementors, such as how to implement sets, structures of well as guidel files, and the standard proceauce dispose, as with other Pascal implementations. 'PASCAL NEWS' has provided this, but the Implementors Group Newsletter could provide extended and more specialized features.
~ Provide a means for implementors to add to a 'validation suite', such as that mentioned in PASCAL NEWS' \#12 (pages 52 \& could contribute a program that gave him fits, and see if developers had solved that problem.
~ Provide an information exchange for all persons interested in Pascal. The PuG offices are inundated with requests for information on a particular implementation, and the Implementors Group could serve as a clearinghouse, to channel appropriate persons.
~ Provide an organization to evaluate and decide on proposals for Pascal extensions. This may include the experimental language features suggested by the UCSD Workshop or 'conventionalized extensions' proposed by other persons.

Provide a forum for implementors and users to interact with each other. This includes user's comments about a particular implementation, to give implementors a better feel for what users think of their compiler (many developers share the compilers, and have no feeling for the number of compilers in use or how successful they have been), as well as the checklists which implementors now provide for 'PASCAL NEWS'.

Tnis list is not complete by any means, and we are looking for suggestions of other areas where the group can serve all pascal users.

The logistics of this group still have to be worked out. I have volunteered to act as the group's coordinator, at least will be open to all interested pascal User's Group mempers, and a mailing list will be maintained so that all known implementors and interested persons receive the newsletters and other mailings. There will be no fees until we have a feeling on the number of people involved, and the cost can be determined.

Everyone who is interested in this group would have to be prepared to contribute in some form. This could include responding to queries regarding their particular implementation, and contributing to the newsletter. I can provide the clearinghouse function and forward inquiries to the right people, but I don't have the time or the knowledge to be able to answer

## Implementation Notes

questions on all compilers. The same is true of the newsletter t is for all implementors, and is an excellent venicle to show ff an elegant solution to a sticky problem, as well as
convenient means to communicate with other Pascal implementors.
Please send any comments,etc. or requests to de on the mailing list to me:

> Scott K. Jameson Hewlett-Packard lloo Nolfe Road Cupertino. CA 95014
(408) 257-7000

## CHECKLIST Pascal Implementations Checklist

0. DATE/VERSION
(* Last checklist changes; version name or number, if any. *)
1. DISTRIBUTOR/IMPLEMENTOR/MAINTAINER
(* Names, addresses, phone numbers. *)
2. MACHINE
(* Manufacturer, model/series and equivalents. *)
3. SYSTEM CONFIGURATION
(* operating system, minimum hardware, etc. *)
4. Distribution
(* cost, magnetic tape formats, etc. *)
5. DOCuMENTATION
(* In form of supplement to Pascal User Manual and Report? In form of supplemen
Machine retrievable? *)
6. MAINTENANCE POLICY
(* How long? Accept bug reports? Future development plans. *)
7. STANDARD
(* Implements full standard? Why not? What is different? *)
8. MEASUREMENTS
(* -compilation speed (in characters/sec. please; this is a meaningful measurement for compilation speed); -compilation space (memory required at compilation time);
-execution speed;
-execution space (the memory required at execution time;
compactness of abject code produced by the compiler);
** Try to compare these measurements to the other language
processors on the machine, e.g., FORTRAN. *)
9. RELIABILITY
(* stability of system (poor, moderate, good, excellent); how many sites are using it?
when was the system first released to these sites? *)
10. Development method
(* Compiler or interpreter? Developed from Pascal-p / handcoded from scratch/bootstrapped/cross-compiled/etc.? What language? Length in source ines? Effort
11. LIBRARY SUPPOR
(* Libraries of subprograms available? Facilities for external and FORTRAN (or other language) procedures available? Easily linked? Separate compilation available? Automatic copy of text from library into source program available? Symbolic dumps available? *)

Return to: Pascal Implementations
c/o Andy Mickel Center: 227 EX
University of Minnesota
Minneapolis, MN 55455 USA

PORTABLEPASCALS

Pascal-p. idealized architecture for Pascal. The symbolic code is thus called P-code.
On the magnetic tape are textfiles containing:

- a sample character set collating sequence. This file is also distributed as a listing to simplify character set conversion
-the Pascal-P compiler in Pascal. write an interpreter in an existing language on the target computer system. - a Pascal-P compiler in P-code. In other words, the result of compiling the Pascal-P compiler on itself.

The person implementing Pascal has several choices. If there is no access to a working Pascal compiler on another machine, the implementor orders a Pascal-P kit already configured are given by the implementor on the Pascal-P order form. (See below.)

After receiving the kit, the implementor can write an interpreter for $P$-code in another language (usually takes about one person-month), and thus immediately has access to a Pascal compiler running interpretively by using the P-code version of the compiler included. in the kit

To produce a real Pascal compiler for the target machine then requires editing of the Pascal-P compiler written in Pascal to produce code for the target machine (instead of the P-machine). After recompiling, a Pascal compiler exists in the code of the target machine.

If the implementor initially has access to a working Pascal compiler on another machine, the step of writing a P-code interpreter can be omitted.

Facts about the Pascal-P compiler:


#### Abstract

The most-widely used portable compiler for creating new Pascal Implementations is Pascal-P. Basically Pascal-P is distributed from three places in the form of a kit Pascal-P. Basically Pascal-P is distributed from three places in the form of a kit consisting of a magnetic tape and printed documentation. Pascal-P is a compiler written in Pascal (almost 4000 lines) which generates symbolic cal-P. Basically Pascal-P is distributed from three places in the form of a kit sisting of a magnetic tape and printed documentation. Pascal-P is a compiler written in Pascal (almost 4000 lines) which generates symbolic code for a hypothetical stack machine called a "P-machine" because it is a low-level


- a P-code assembler/interpreter written in Pascal which is intended to document how to
- The current version is called Pascal-P4 and is distributed with a copy of Pascal-P3 (which is of interest to previous recipients of Pascal-P2).
- Pascal-P4 represents a major improvement over earlier Pascal-P versions because it removes data-type-alignment restrictions, is more efficient, includes runtime tests, and is a more complete implementation of Pascal.
- Pascal-P2 was developed from a phase in the stepwise refinement of Urs Ammann's Pascal-6000 compiler in 1974 by K. V. Nori, Urs Ammann, K. Jensen, and H. H. Nageli. Subsequent improvements were done by Christian Jacobi
- Reliability of Pascal-P4 has been fairly good. As of Spring, 1977, it was distributed to 106 sites by George Richmond (from Colorado), to 37 sites by Chris Jacobi (from Switzerland), and to more than a dozen sites by Carroll Morgan (from Australia).
- The is no promise of maintenance for Pascal-P. P4 is the final version produced at Zuerich. We do print reports of bugs (and fixes) in P4. Over 25 fixes were printed last year in Pascal News issues \#11 (pp 70-71) and \#12 (pp 56-57). More are printed below
Documentation for Pascal-P4 consists of a 65-page report entitled The Pascal $\langle P>$ Compiler: Imp.ementation $\frac{\text { Notes }}{\text { to the original December, }} 1974$, edition is also available.)
Pascal-p ${ }^{\text {does not adhere strictly to Standard Pascal (th }}$
(o Standard Pascal (the User Manual and Report).

1. nil is implemented as a predeclared constant, and forward as a reserved word. The standard indicates that nil is a reserved word, and forward is not listed as a reserved word.
2. The standard comment delimiters \{ and \} are not supported.
3. The following standard predeclared identifiers are not provided: maxint, text, found, page, and dispose. Further, the following standard predeclared identifiers are recognized but are flagged as errors: reset, rewrite, pack, and unpack.
4. The program heading is not required by $P 4$
5. Non-discriminated variant records are not supported.
6. The compiler does not allow a ";" before the "end" in a record type. (See the P4 bug reports in Pascal News \#12 (pp 56-57) for a fix.)
7. None of the following file-related features are supported
-- Declaration of file types, variables, and parameters
解
-- The requirement by the standard that the standard files input and output appear in the program heading if they are used
-- Access to non-text files using read and write. form with write.
8. Formal-procedures and formal-functions are not supported.
. Set constructors containing the subrange notation (e.g., ['0 $\left.0^{\prime} .^{\prime} 9^{\prime}\right]$ ) are not supported.
9. "Non-local" goto statements are not supported.

Pascal-P can be ordered from three places (write for prices and order forms)
In Europe, Asia, and Africa, order from:
hristian Jacob
nstitut fuer Informatik
CH-8092 Zentrum
CH-8092 Zue
Phone: 41/1-32 $6211 \times 2217$

In North and South America, order from:

In Australasia order from
ascal Distribution
o Steve Winograd ${ }^{6}$.
University of Colorado
Boulder, CO 80309
Boul
USA
Phone: 303/492-8131
Tony Gerber
Basser Dept. of Computer Science
niversity of Sydney
Australia
hone: 6I/2-692 3216

Pascal P4 -- Bug Reports

78/06/09, Ted C. Park, Systems Development Medical Data Consultants, 1894 Commercenter West - Suite 302, San Bernardino, CA 92408 (714) 825-2683, reported:
I just came across two more bugs in the PASCAL-P4 compiler. FUNCTION EQUALBOUNDS contains an obvious error
replace P. 136 with GETBOUNDS(FSP2,LMIN2,LMAX2);
ROCEDURE GEN2T is used for (among other things) generating 'CHK' instructions. The fix causes the width of the 'p' field to be 3 or 8 as needed. Without the fix the lower limits of arrays must be less than four digits long!
replace P. 262 with $\operatorname{WRITELN}(\operatorname{PRR}, \mathrm{FP1}: 3+\operatorname{ORD}(\mathrm{ABS}(\mathrm{FP1})>99) * 5, \mathrm{FP} 2: 8)$; \#9-10 (p 63) and \#11 (p 72).
(*Thanks Ted!*)
Pascal Trunk Compiler

The trunk compiler is the machine-independent part (e.g., syntax analysis and error recovery) of a Pascal compiler in which the code generation has to be inserted in a certain number of empty procedures. We have received no new information on the
compiler since that which we published last year in Pascal News issue $\# 9-10$ ( p 62 ).

Pascal-J is a compiler which translates Pascal to the intermediate language Janus, a totally portable "mobile programming system" -- even to the point of defining its own character set! Janus in turn is macro-processed via Stage2 which is implemented in standard Fortran. We have received no new information on Pascal-J since that which we published last year in Pascal News issue \#9-10 ( $p-62$ ).

## Pascalvariants

Pascal-S
Pascal-S is a subset pascal developed by Niklaus Wirth. We have received no new information on Pascal-S since that which we published last year in Pascal News issues
Pascal J

A portable pair of Pascal compilers was implemented by Per Brinch Hansen and Al Hartmann at Cal Tech in 1974-1975 for the PDP 11/45. The system consists of a "Sequential Pascal" compiler, a "Concurrent Pascal" compiler (used for writing operating systems and other concurrent programs), and a "kernel" or machine dependent set of run-time routines
written in assembler. The project at Cal Tech centered around writing a one-user operating system called SOLO in Concurrent Pascal. Both compilers are written in Sequential Pascal

In 1975-1976 the system was distributed widely ( 252 sites) and led to the development of a machine-independent version with a different kernel.

The distribution tapes ( $\$ 50$ ) and documentation ( $\$ 10$ ) can be ordered from:
Pascal Distribution
c/o Steve Winograd
c/o Steve Winograd
Computing Center: 36
-University of Colorado
USA
USA
Phone: 303/492-8131
Publications about Concurrent Pascal include:
(1) "The programming language Concurrent Pascal", in the June, 1975, IEEE
(2) Transactions on Software Engineering $1: 2$, by Brinch Hansen.
(2) A guest editorial and four articles by Brinch Hansen in the April-June, 1976 issue of Software - Practice and Experience 6, pp 139-205. The articles are entitled:
The Solo Operating System: A Concurrent Program"
"The Solo Operating System: Procedures, Monitors, and Classes"
"Disk Scheduling at Compile Time"
(3) The book Operating Systems Principles by Per Brinch Hansen, Prentice Hall, 1973.
(4) An article "Experience with Modular Concurrent Programming" in the March, 1977,
(5) IEEE Transactions on Software Engineering 3:2, by Brinch Hansen.
(5) A Concurrent Pascal Compiler for Minicomputers by Al Hartmann, Springer-Verlag:

Lecture Notes in Computer Science, Volume 50, 1977.
(6) The new book The Architecture of Concurrent Programs by Brinch Hansen,

## COMPUTER SCIENCE DEPARTMENT

Salvatori Computer Science Center
(213) 741-5501

October 1, 1978

Dear Concurrent Pascal User,
It is now 3 years ago since the Concurrent Pascal compiler and the Solo Operating System were first distributed. Since then the system has been moved to several computers and used for a variety of purposes.

Some users (but not all) have briefly reported on their usage of Concurrent Pascal in the Pascal Newsletter. I am now trying to get a more complete overview of the current use of the system.
If you are using Concurrent Pascal or Solo then please send me a letter. I would like to know which computer you are using, how hard it was to move the system to that machine, how reliable the and any other comments you may have. I would also like to know i you have published any papers about your experience.

Andy Mickel and I plan to publish these letters in the Pascal Newrs If I receive your letter before February. 28, 1979 it will be included in the newsletter.

I look forward to hearing from you.



## UMIST

The University of Manchester Institute of Science and Technology PO Box 88, Manchester M60 10D Telephone 061-236 3311 Department of Computation

27th April 1978

Dear Andy,
We have moved Brinch Hansen's SOLO system on to our 40 K CrIL Modular One computer. We have found the system to be very reliable and the few bugs that have been found have been simple to fix. Gur main interest is in Concurrent fascal which we are using as a tool for our work on the development
of programming methods for multiprograms.

Due to the inhospitable arcinitecture of the Modular One our system runs at only a fifth of the $s_{:}$eed of the original PDP-11/45 implementation. iork is under way to improve the speed by the utilisation of a second processor and a fixed head disc. $\dot{A}$ simple multi-acce:is system is also being considered.

The transportation of SOLO was very straight forward and was accomplished in about eight months by two undergraduate students and one lecturer working part-time. Further details of the move are contained in Fialcolm fowell's report [1].

We are interested in exchanging information and programs with other users or potential users of SOLO or Concurrent Pascal.

Yours sincerely,
Derch Colenar
Derek Coleman
Lecturer in Computation
[1] M.S. Fowell. Experience of Hovine and Usint the SULC Operating System, Comoutation Department, Uilis

## Modula

Modula is a small language for dedicated computer systems and process control applications on small machines, developed by Niklaus Wirth and co-workers in 1975-76. It is conceptually cleaner than Concurrent Pascal in many respects. The Modula language definition provides for machine-dependent facilities for interacting with asynchronous devices. Modula is still experimental and the implementors in Zurich have insisted there are no distribution arrangements. Other implementations are complete or underway. See
Pascal News $\# 11$ (p 74) for details of the University of York PDP-11 compiler. Also, on $78 / 10 / 27$, Gerd Blanke (Postbox 5107; D-6236 ESCHBORN Germany; phone (06198) 32448) wrote "MODULA will be running on a ZILOG MCS with 64 K under RIO near the end of this year!"

Published material on Modula includes:
(1) "Modula: A Language for Modular Multiprogramming", Software - Practice and Experience 7 (1977), pages 3-35, by Niklaus Wirth
2) "The Use of Modula", same as (1), pages 37-65, by Niklaus Wirth.
(1), pages $67-84$, by Niklaus
4) Wirth.
(August a Discipline of Real-Time Programming", Communications of the ACM 20:8
5) "Experience with the programming language MODULA", University of York - Dept. of Computer Science (June, 1977), by J. Holden and I. C. Wand
References (1) through (3) received very interesting reviews in Computing Reviews 18 (November, 1977), \#32217, \#32218, and \#32219.

## Feature Implementation Notes

## IMPLEMENTATION NOTE

Implementation of INPUT and OUTPUT Arthur Sale and Judy Bishop

## PROBLEM

It has come to our attention that there is a problem with the implementation of the pre-defined files input and output. What follows refers only to output, it is easier to demonstrate the effects on an output file, but applies equally to the file input.

The problem turns on two of Pascal's Achilles' heels: the elision of a file-name in read and write and the resulting default, and the singuiar program parameter part and its interaction with pre-defined names. The situation can be summed up by two questions, to each of which there are two reasonable answers.
QUESTION 1 : Where do default writes go?
Does write $(x)$ write on the default file named output (and pre-defined), or
on the lexically innermost definition of a file named output?
Answer A : x is always printed on the pre-defined file, whatever
nswer B : the symbol table is searched for output and the write i attempted on the innermost occurrence of it.

## QUESTION 2 : At what level is output defined?

Is the pre-defined file output regarded as declared at the level of the
program block (level 0 ) or in a lexically enclosing block (level -1 )?
Answer C : the file is regarded as being at the level of the
the file is regarded as being at the level of the program that level.
Answer D : the file is regarded as being in a block enclosing the program, so that the name can be redefined in the program block.

HERE DO DEFAULT WRITES GO?
The Tasmania B6700 compiler and the AAEC IBM compiler transmit default information always to the pre-defined file output, and it seems likely that the CDC-6000 compiler, the ICL compilers, and most Pascal-P derivatives do the same. These indicate that Answer A is currently predominant.

What does the Report and User Manual say? The Report (\#12.3) defines write (x)
as equivalent to
write (output, $x$ )
hich makes one think of Answer B: the elision of the file-name is to be handled by a macro-expansion. However, on reading the User Manual (p61) and parameter which is assumed by default if the filename is omitted. $\frac{\text { In other }}{}$ words: the pre-defined file, and Answer A.

On balance, therefore, the predominant Answer A seems to be approved by the User Manual. It can be argued that this is abstractly best, for if we have oo have any defaults in Pascal (and we've got these few), then they ought to be as simple as possible.

The following is a test program to exercise your compiler and test its performance on this question:
program question1 (output);

## procedure inner

var output : text;

## begin

writeln('WRITING ON DEFAULT FILE')
writeln(output,'WRITING ON LOCAL FILE')

## end;

begin
writeln(output,'TEST OF QUESTION 1')
inner;
writeln('RAN')
end.
at what level is output defined?
output, in common with other pre-defined names, can be regarded as pre-declared in a lexically enclosing. scope, thus allowing its redefinition in the program block. This is asserted by the Report (p161), and is the current interpretation
given by the Tasmania B6700 compiler: Answer D.

The alternative, sanctioned by the User Manual (p91) in the CDC-specific section, says that these files are implicitly declared in the program block (not pre-declared). In CDC-6000 Pascal therefore one may not define any object with the name output at the program level. The AAEC compiler is similar, thus giving Answer C.
So both answers find some support, and both are in use. Which is better? Experience of one of us (AHJS) indicates that perhaps Answer $C$ is best: pre-declared at the program level. This experience arises from a number of apprived from a user atto refte the output file by declaring a hidin occurrence of the name output at the program level. If Answer C had been adopted in the Tasmania 66700 compiler, these would have been detected as illegal by the compiler, and other name choices would have been forced on the users.
tis also possible, but inconclusive, to argue from analogy. Focussing on the analogy with other pre-defined identifiers, such as abs and true, then it program block. But, of course, these two files (input and output) are the only two var objects which are pre-defined, so perhaps they should be special. This view leading to implicit declaration in the program block, is supported by the analogy with all other file names mentioned in the program parameter part which must have a declaration in the program block (at least in the CDC-6000 implementation; others allow more freedom).
If then Answer C is more attractive, the Tasmania 86700 compiler should be changed. In this case however, we shall wait until the draft standard for ascal resolves the issue. The following test program will show what your compiler does:

```
program question2(output)
    var output : integer;
    { if this compiles, you've probab&y got Answer D
begin
    output := 1;
end.
```


## DEVIANT IMPLEMENTATION

ne implementation, which shall remain nameless as a fitting punnishment, lies outside the permitted limits of the Report and User Manual by using a subtle hange. In this implementation, elision of the file-name causes the write to take place on an un-named pre-defined file. This has the result that write (output, x)
fails to compile unless another file is declared with the name output, and that the question of the default file's scope does not arise (because you an rename to confusion amongst users.

## BRIEF ADVICE

## To Pascal users:

(1) Do not use the identifiers input or output for anything other than the pre-defined files that you don't need to declare.
2) Preferably do not leave the file-names out of reads and writes, but put them in explicitly as a good programming practice.

## To implementors

please modify any implementation plans to be consistent with majority opinion in Answer A, and watch for more information on Question 2

## To language designers:

(1) Future languages should make it mandatory for compilers to inform users of any names they hide under scope rules, if such exist. The extended searches are only necessary at declaration points.
(2) Defaults of any kind should be avoided

Arthur Sale
1978 June 13
(Revised-1978 August 1)
University of Tasmania
Judy Bishop
University of the Witwatersrand

1978 April

## IMPROVED CHECKING OF COMMENTS

As is well-known; comments of the PASCAL kind have a severe disadvantage in that if a closing marker is omitted or mis-keyed, intervening source text will be treated as commentary until a closing marker is found for a later comment. Since such errors do not give rise to syntax errors, they may remain undetected in source text for a long time. This feature is exacerbated in PASCAL by allowing good compile-time error-detection.
In Burroughs B6700/B7700 PASCAL (University of Tasmania compiler), the problem this creates for programmers (especially learners) has been alleviated by issuin be the result of an error. Very few within a comment, as remain undetected, and the change in the lexical analyser is very simple. This suggestion is commended to other implementors and maintainers.

People who use the comment facility to suppress source text compilation (debug code; superceded text) may be annoyed by the many warning messages. They can then suppressed by our compiler option WARNINGS; but better still would be to realise that this is a misuse of comments and hardly likely to enhance readability!

Many Algol 60 compilers have included similar checks in their handling of the singularly nasty end-comment in that language; the experience is generalizable to PASCAL too.

It would be possible to issue warnings for other symbols encountered in comments for example a comment opening marker, and this would marginally improve the detection probability of these errors. We judged such extension as not worth the effort, especially since both \{\} and ( $* *$ ) comments are permitted in our PASCAL,
which would require quite complex checks.

## Sample output for error:

\{this comment is unclosed
count := 0;

WARNING: DISCOVERED ";" IN COMMENT. DID YOU FORGET TO CLOSE A COMMENT
\{this closing marker will match the first one\}

Lazy I/O
(* The "Lazy I/O" scheme has apparently been invented several times. The earlies mill discussed in some detail at the UCSD Workshop in July. The consensus there seemed to b that Lazy I/O is the best solution anyone has yet proposed, even though it may be somewhat ess efficient than other approaches in terms of execution time,
On October 21, James Saxe and Andy Hisgen added a note written to Andy Micke1 which said: "By the way, the lazy evaluation idea was not cribbed from Berkeley UNIX Pascal, as you have suggested, but was developed here independently. We are, however, glad to see that there are other peple around who do not "1 Jin the semantics of Pascal at every opportunity." - Jim Miner *)

> Comioüäter Sciēāe Dedärtrient Carnegie-Mellon University Pittsburgh, PA 15213 August 4, 1978

Subject: LAZY EVALUATION OF THE FILE BUFFER FOR INTERACTIVE I/O

## Dear Andy,

A frequently occurring difficulty in Pascal programming, and one which is particularly puzzling to the novice, arises from the effect of the file lookahead buffer on interactive I/O. Specifically, let TTY be a TEXT (FILE OF CHAR) variable associated with the input stream from the user's terminal and let TIYOUPUT be a TEXT variable associated with the output stream to the terminal. Now, consider the following program fragment:

$$
\begin{aligned}
& \text { ReadLn (tty, nplayers); } \\
& \text { WriteLn (ttyoutput, Number of marbles = ?'); } \\
& \text { ReadLn (tty, nmarbles); } \\
& \text {. . . }
\end{aligned}
$$

Under many Pascal implementations, this fragment will fail to work as intended because the READLN in line 1 will not complete until the lookahead buffer, TTY + , has been filled with a character (presumably the first digit of NMARBLES) from the terminal. The user, meanwhile, will not supply this character until he has been prompted by line 2, which of course cannot happen until line 1 has finished execution. Attempted "solutions" to this problem include

- Use of special user-defined procedures for the terminal which read a single real or integer after doing a READLN. (This approach, of course, is not very useful for programs that do character input.)
- Altering, in various ways, the semantics of file input when the run-time system "knows" that the file being read happens to really be the terminal (e.g., making EOLN(TTY) be FALSE and TTY' be '' after each READLN(TTY) regardless of the contents of the following line of input. Note that empty input lines will no longer be reliably detected and may "hang" the terminal.).
- Introducing a new file type for interactive devices, with slightly different semantics from those for TEXT files.

We maintain that all these. kludges are completely unnecessary. A Pascal compiler and run-time system can be made to support interactive I/O in a perfectly natural manner without any deviation from the semantics laid out in the report. This can be achieved by "lazy evaluation" of the file lookahead buffer for the terminal, that is, the practice of never filling TTYY until it is actually used.

To describe this more precisely, let ACTUALGET be a procedure having the effect that GET has in most implementations. That is,

## ActualGet (tty)

has the effect of grabbing one character of terminal input from the operating system and sticking that character in TTY. We introduce a new Boolean variable, TTYBFULL, visible only to the run-time system, which, as we shall see, shall be TRUE iff the "current" character in the file TTY has actually been read from the
terminal. The action of
Get ( $t$ ty) ;
is now precisely defined as
IF ttybfull
THEN ttybfull := FALSE
ELSE ActualGet (tty);
Whenever the programmer explicitly does something that requires lookahead (assigns to TTY $\dagger$, calls EOLN(TTY) or EOF(TTY), uses TTY + in an expyession, or (assigns to TTYt, calls EOLN(TTY) or EOF TTY+ as a value parameter), the run-time system, behind the programmer's back, forces the lookahead buffer full by doing

> IF NOT ttybfull THEN
> BEGIN
> ActualGet (tty)
> ttybfull : $=$ TRUE;

END;
When TTY is RESET for input form the terminal an ACTUALGET is not done, but TTYBFULL is initialized to FALSE. The call

Read (tty, c); (* where C is a variable of type CHAR *)
continues to be equivalent (as specified in the Report) to

> c := tty个;

Get (tty);

The procedure ACTUALGET, like the variable TTYBFULL, is directly accessible only to the run-time system and not to the programmer.

Careful consideration of the rules described above will show that they result in exactly the semantics described in the Pascal Report. The only difference between this and other implementations is that the terminal will not "hang" in the manner described in the opening paragraph. This conformity with the semantics of the Report has several advantages:

- Conformity to Standard Pascal improves the prospects for software portability.
- Any program which works correctly under a correct implementation of Standard Pascal will continue to work, and will give the same output (given the same input), under the implementation described above.
- Since the semantics of disk file I/O and terminal I/O continue to be identical, programs which use input from one source can be easily modified (say, for debugging) to take input from the other. Also programs which postpone until run-time the decision whether to take input from a disk file or from the terminal can be written without needless duplication of code.

Let us emphasize again that even programs which make use of the lookahead buffer will work in the manner defined by the Report, because any program action which actually requires knowledge of the lookahead character will demand that character from the terminal before it can continue. Of course it is the programmer's responsibility to prompt the user for this input, but since the programmer knows that this information is required at a particular point in the program, he should have no trouble remembering to prompt for it Consider, for example, the following program fragment, which prompts the user for an integer but allows him to just type a carriage return if he wants the default value (shown in brackets by the program):

```
WriteLn (tty, 'Number of runs [10] : ');
IF Eoln (tty) THEN
    BEGIN
    nruns := 10;
    ReadLn (tty);
ELSE
ReadLn (tty, nruns);
```

. . .

In line 2, the programmer does an explicit lookahead at the first character on the line to determine whether it is the line delimiter (i.e., whether the line is empty). In this case, the lookahead character will be demanded by the run-time system before the expression EOLN(TTY) can be evaluated. However, the prompt for this input will have already been supplied by line 1.

In closing, we should take note of some tricky aspects of the lazy evaluation technique which might at first escape the notice of the prospective implementor. First, lazy evaluation of the lookahead buffer should be performed on all TEXT files, since it is not necessarily possible to determine at compilation which of these will be associated with the terminal (for example, TTY may be passed as an actual procedure parameter). Second, enforcing correct semantics can be very tricky in procedure parameter). Second, enforcing correct semantics can be very tricky in
cases where the lookahead buffer (TTYt) is passed as a VAR parameter [Our cases where the lookahead buffer (TTYt) is passed as a VAR parameter [Our
approach at CMU is to force the buffer full once at the time of function or approach at CMU is to force the buffer full once at the time of function or procedure invocation and to leave the user on his own thereafter. Since passing
TTY rather than TTY guarantees the expected semantics, we feel that this approach does not make impositions on the reasonable user. An alternative approach does not make impositions on the reasonable user. An alternative approach would be to disable lazy evaluation for the duration of the invocation.]. In spite of these difficulties, however, we believe that the lazy other mechanisms we have seen.

Yours truly,


Lindy Hicoun
Andy Hisgen

## Alpha Microsystems AM-11

See DEC LSI-11 UCSD.

Altair 680b
See Motorola 6800 St. Paul.

Altair 8800

See Intel 8080.

Altos ACS-8000
It has been reported that Altos Computer Systems; 2378b Walsh Ave.; Santa Clara, CA 95050; 408/244-5766 offers a Zilog Z-80 based microcomputer which supports CP/M an ascal, but we have received no information from Altos

Amdahl 470

See also IBM 360/370.
It has been reported that the IBM $360 / 370$ AAEC as well as the Vancouver systems are running on an Amdahl 470.

Andromeda 11/B

See DEC LSI-11.

## Apple II

See MOS Technology 6502.

BESM-6 Moscow
0. DATE/VERSION. 78/9/21.

1. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER. S. Pirin; Moscow Computer Center; USSR Academi Sciences; Moscow, R.S.S.R.; U.S.S.R.; (* No phone number reported *)
2. MACHINE. BESM-6.
3. SYSTEM CONFIGURATION. (* No information reported. *)
4. DISTRIBUTION. (* No information reported. *)
5. DOCUMENTATION. (* No information reported. *)
6. MAINTENANCE. (* No information reported. *)
7. STANDARD. (* No information provided. *)
8. MEASUREMENTS. (* No information reported. *)
9. RELIABILITY. (* No information reported. *)
10. DEVELOPMENT METHOD. (* Reported that project has been underway (or possibly complete?) for some time. *)
11. LIBRARY SUPPORT. (* No information reported. *) BTI 8000

It has been reported that the BTI 8000, a 32 bit multiprocessor system offered by BTI Computer Systems; 870 W Maude Ave.; Sunnyvale, CA 94086 ; 408/733-1122, includes a Pascal compiler bundled with the hardware and that the system software is written in "Pascal-X", Burroughs B1700 Zurich

We have received no new information on this implementation since that which we published last year in Pascal News issues: \#9-10: 73. \#12: 57-58.

Burroughs B1800

See Burroughs B1700 Zurich.
Burroughs B4700 Fredonia

We have received no new information on this implementation since that which we published last year in Pascal News issue: \#9-10: 73.

Burroughs B5700 Edinburgh


3 May 1978

## EE: PASCAL for Burroughs B5700

This compiler is in current use here and is available from me. Imagine it's also available from University of Wisconsin - Eau Claire as noted in the December, 1977 "Pascal Answers". The original source of the compiler is Heriot-Watt University, Edinburgh; and any complaints, bugs, fixes, etc. should be sent there.

## No charge if the requester sends a tape



James H. Haynes
Associate Development Engineer

## Burroughs B6700 Helsinki

According to Antti Salava (* 78/10/18 *) : "I'm not working with Pascal nowadays. A year ago I left the University of Helsinki, where I was implementing Pascal-HB compiler on the Burroughs B6700. It's been running now a couple of years without any fatal crashes. We wrote a report on our compiler, too. Hasn't anybody noticed it? It's this: Hannu Erkio, Jorma Sajaniemi, Antti Salava; "An Implementation of Pascal on the Burroughs B6700"; Department of Computer Science; University of Helsinki; Report A-1977-1. Copies may be
ordered from: Department of Computer Science; University of Helsinki; Toolonkatu 11; SF-00330 Helsinki 10, Finland.

## Burroughs B6700 San Diego

We have received no new information on this implementation since that which we
 a copy of the multi-page machine-retrievable installation notes that come with the system.

Burroughs B6700/7700 (Tasmania)
0. DATE/VERSION. Checklist has not been updated since 78/03

1. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER. A.H.J. Sale; Pascal Support; Dept. of Information Science; University of Tasmania; Box 252C G.P.O.; Hobart, Tasmania 7001 Australia; STD 002 23-0561 x435.
2. MACHINE. Burroughs Model III B6700, B7700.
3. SYSTEM CONFIGURATION. Burroughs MCP version II. 8 (with few (minor) local mods). Minimal system to operate not known, but unlikely to be any B6700 that small--storage demands are low, and little else is critical.
4. DISTRIBUTION. Both 7- and 9- track magnetic tapes available. Annual fee of \$100 (Australian) is charged to cover mailing, processing, and maintance costs, payable to "The University of Tasmania".
5. DOCUMENTATION. Available documentation: Report R77-1: Supplement to Pascal User Manual and Report ; Report R77-3: Reference Manual similar to B6700 ALGOL's; A Pascal Language *)
6. MAINTENANCE. To be maintained for teaching use within the University as well as larger aims. Reported bugs will be fixed as soon as possible, with patch notices to users Duration of support not yet determined; several other developments are also pending. Each installation will be issued a supply of FTR-forms similar to those used by Burroughs for use in corresponding with us, and we will attempt to do a professional job in maintenance of the system.

The compiler has been stable in code for some time, reflecting its basic integrity. new versieatures are added from time to time, and notified to users as patches or as a new version release. The department accepts FTR notices, and will attempt to fix those
which warrant such attention．Some modifications have taken place as a result of user feedback．The compiler was especially designed so as not to generate dangerous code to the
MCP，and no system crashes have been attributed to it since the first few months of MCP，and no system crashes have been attributed to it since the first few months of testing，and then only three．
7．STANDARD．
Restrictions：Program heading：reserved word program is synonymous with procedure；no parameters（files）are permitted after the program heading．Reason：CDC anachronism of no utility in our installation，and likely to be confusing．Set constructor of form A．．B not Reason：a ridiculous feature to standardize．Full Pascal I／0 not implemented．Reason： future plans．Present I／scheme is like Pascal－1．

Extensions：otherwise in case statement．Various reserved words，character set transliterations．Burroughs comment facility．File attributes in declaration．Format declarations．Extensive Burroughs－compatible compiler options．（Pascal control comment option mode not implemented）．

## 8．MEASUREMENTS．

compiles about $20 \%$ slower than FORTRAN or ALGOL，but in about $2 / 3$ of their space（for test programs about $4-5 \mathrm{~K}$ words on average instead of $8-10 \mathrm{~K})$ ．Elapsed compilation times similar，though pascal slower．Speed should be improved by eventual tuning．
executes at same speed as FORTRAN and ALGOL（code is very similar and optimal）and takes generally longer elapsed residence time primarily due to MCP intervention to create new segments for residence times about $20 \%$ greater than equivalent ALGOL．

9．RELIABILITY．Excellent．Only one system crash during testing attributed to pascal． Compiler now in use at 3 sites．True compiler has been in use since 76／10．First released to outside sites in 77／4．
10．DEVELOPMENT METHOD．Compiler which generates $B 6700$ code－files which are directly executed by the B6700 with MCP．Written entirely in B6700 ALGOL．Hand－coded using Pascal－
as a guide／model．All other paths offered much more difficulty due to special nature of machine／system．Person－month details not kept，and project proceeds in fits and starts as teaching intervenes．Project has thus far been limited to two people：Prof．A．H．J．Sale and R．A．Freak（Support programmer）．

11．LIBRARY SUPPORT．There is as yet no BINDINFO in the code－file so that it is not possible to link Pascal to modules compiled by other language processors，but the system contains an extended set of predefined mathematical functions

CDC 2550

See CDC Cyber 18 La Jolla

CDC 6000，Cyber 70， 170 Bethlethem，PA

We have received no new information on this implementation since that which we published last year in Pascal News issue：$⿰ ⿰ 三 丨 ⿰ 丨 三$ 11： 82.

CDC 6000，Cyber 70，Cyber 170 （Zurich）

0．DATE／VERSION．Pascal 6000 Release 3；78／11／15．

Wally Wedel
Computation Center
Switzerland
University of Texas－Aus tin
Austin，TX 78712 USA
Austin，TX 78712 USA
$512 / 472-3242$
512／472－3242
－（Australia，New Zealand，or Oceania） Basser Dept
Basser Dept．of Computer Science Sydney，N．S．W． 2006 Australia
61 ／2－692 3216
612／376－7290
＊Arrangements are underway to have the implementor of the CDC 7600，Cyber 176 run－time system take over distribution for Europe，Asia，and Africa from the original implementor．

2．MACHINE．Control Data 6000 series，Cyber 70 series，and Cyber 170 series．
3．SYSTEM CONFIGURATION．Minimum central memory－49K words．Operates under Scope 3．4，
4．DISTRIBUTION．Tape format is Scope internal binary 7／9track，unlabelled， 800 bpi ． Specify：person responsible for maintaining the system，your hardware，operating system， machine－retrievable source and object decks，installation notes，and software tools． Arrangements for distribution（cost，etc．）for the new release have not yet been finalized．Contact the distributor in your area in further information．
5．DOCUMENTATION．Machine－retrievable supplement to Pascal User Manual and Report． Documentation of library－support package is available with Release 3.

6．MAINTENANCE．Will accept bug reports at Minnesota for forseeable future．
7．STANDARD．Nearly full standard．［Restrictions include：standard procedures and functions cannot be passed as actual parameters；file of file is not allowed］［Extensions include：additional predefined procedures and functions；segmented files，conformant array parameters，otherwise in case statement，variable initílization facility（value），and text－inclusion facility for source libraries．］
8．MEASUREMENTS．
Compilation speed：10800／5800 characters per second on a Cyber 74／Cyber 172；
Compilation size： 40 K （octal）words for small programs； 57 K for self－compilation．
Execution speed：seif－compiles in $65 / 120$ seconds．
Execution size：binaries can be as small as 2.4 K ，compared with Fortran minimum of over 10K．
9．RELIABILITY．Unknown，as this is a new release．However，Release 2 was very reliable and was in use at over 300 knowin sites．First version of this comipiler was operational in late 1970．The present version was first released in May 1974．A pre－release version of release 3 was tested at 10 sites for up to 5 months prior to the official release．

10．DEVELOPMENT METHOD．Bootstrapped from the original Pascal－6000 compiler，but developed in a 6－phase stepwise－ref
rewritten for Release 3 ．

11．LIBRARY SUPPORT．Allows calls to external Pascal and assembler subprograms and Fortran （FTN）subroutines．The user library supplied with the system contains many intrinsic procedures and functions in addition to the Standard Pascal ones．
1．IMPLEMENTOR／DISTRIBUTOR／MAINTAINER．
Distributors：
－（Europe，Asia，or Africa）
See Ric Collins
Univ．of Manchester（CDC 7600）＊
－（North or South America）
mplementor：
Institut fur Informatik
E．T．H．－Zentrum
CH－8092 Zurich

## aintainer：

John P．Strait／Andy Mickel University Computer Center
227 Ex
University of Minnesota rinneapolis，MN 55455 USA
－ 1

0．DATE／VERSION．Release 3 of the CDC 6000 Zurich compiler（from the Minnesota maintainer）is a common release for the CDC 6000，7600，Cyber 70， 170 series．See the letter under CDC 6000 Zurich．

1．IMPLEMENTOR／DISTRIBUTOR／MAINTAINER．This compiler is essentially the Pascal 6000 compiler modified to fit the 7600 and Cyber 76 machines．A new run－time system is being
developed using conditional assemblu in the new Release 3 run－time system by A．P．Hayes； UMRCC；Oxford Road；Manchester M13 9PL；England，U．K．；（061－273 8252）．
2．MAChine．Control Data $7600 \&$ Cyber 76.
3．SYSTEM CONFIGURATION．SCOPE 2.1 .3 or $2.14,32 \mathrm{~K} \mathrm{SCM}$ ．
4．DISTRIBUTION．Contact R．J．Collins at address above．A distribution agreement must be signed and the cost is 30 pounds sterling．

5．DOCUMENTATION．Same as Pascal－6000．
6．MAINTENANCE．UMRCC will assist with bugs－－in the 7600 dependant code（runtime system）only．Minnesota will accept bug reports on the compiler itself．

7．STANDARD．Same as Pascal 6000.
8．MEASUREMENTS．None yet for Release 3；［Release 2 was：Compilation speed is about 57,000 characters $/ \mathrm{sec}$ ．Comp iler comp iles itself in less than 10 sec．Pascal execution speed has been measured by using the obvious encoding in Pascal of Wichmann＇s Synthetic Benchmark（see Computer Journal VoI．19，\＃1）．The Units are in kilo Whetstones．

| compiler and <br> optimisation level | no runtime <br> checking | array bound <br> checking |
| :--- | :---: | :---: |
| ALGOL 4（OPT＝5） | 1996 | 1230 |
| Pascal | 6850 | $6240 *$ |
| FTN（OPT＝2） | 945 | $3174^{* *}$ |

$*$ Using T＋option－－all run time checks included．
$* *$ Forces $O P T=0$ ．
Compiler will recompile itself on a＇half－size＇（ 32 K SCM）machine．Execution space－－Core requirements（octal）： $42,402 \mathrm{SCM}$ ，or 36,045 if segment loaded（using a simple segment structure）．Self compiles in less than 60,000 ．（＊No information provided on size of
compiler or object code produced．＊）］

9．Reliability．Same situation as Pascal 6000 （Zurich）．
10．DEVELOPMENT METHOD．Cross compiled from Cyber 72 compiler．Based on Zurich 6000 compiler with necessary additions for this machine．（＊Person－hours to develop system not reported．＊）

11．LIBRARY SUPPORT．Same as Pascial 6000.
CDC Cyber 18 La Jolla
We have received no new information on this implementation since that which we published last year in Pascal News issues：\＃9－10：75．\＃11：81．

## CDC Cyber 18 Berlin

We have received no new information on this implementation since that which we pubilished last year in Pascal News issue：\＃11：81－82．

7 Tokyo．）

CDC Omega 480－I，480－II

See IBM 360／370．

## CDC STAR－100（Cyber 203）Virginia

We have received no new information on this implementation since that which we published last year in Pascal News issue：\＃9－10： 77.

CII 10070 France

We have received no new information on this implementation since that which we published last year in Pascal News issues：\＃9－10：77－78．\＃12：59－60．（see also Xerox Sigma

GII IRIS 50 Nice

We have received no new information on this implementation since that which we published last year in Pascal News issue：$⿰ ⿰ 三 丨 ⿰ 丨 三 9-10: 77$.

CII IRIS 80 Paris，France
We have received no new information on this implementation since that which we published last year in Pascal News issues：\＃9－10：77－78．\＃12：59－60．（see also Xerox Sigma
7 Tokyo．）

Commodore Pet 2001
See MOS Technology 6502.

Computer Automation LSI－2 and LSI－4 Irvine
We have received no new information on this implementation since that which we published last year in Pascal News issues：\＃9－10：78．\＃12： 60.
CRAY－1 Los Alamos

We have received no new information on this implementation since that which published last year in Pascal News issues：\＃9－10：78－79．

See Zillog Z－80．

27 April 1978

Dear Mr. Mickel:

I am writing to you because of the article that appeared in Computerworld, April 24, on the growth of Pascal and Pascal User's Groups.

The User's Group at Data General will soon have a Pascal Special Interest Group. It is being organized by a member of PUG, Rodney Thayer. He has agreed to serve as an interim co-chairman until the group can elect officers. There will be a Pascal session at the 1978 Annual. User's Group Meeting.
The version of Pascal that we are using is one that has been supplied by R.E. Berry at the University of Lancaster. If you please Eeel free to contact either Rodney Thayer or myself.

Sincerely,

- Clr-wneth. Chegloy
Kenneth A. Roy
D.G. User!s Group
D.G. User!s Group

Richard E. Adams; 967 Atlantic Ave.; Apt. 634; Columbus OH 43229; 614/436-3206 asked (* 78/7/31 *): "I have not seen any references to a Pascal compiler running under Data General's Advanced Operating System (AOS). Is anyone out there working on it?"

Data General Eclipse/Nova Columbia
0. DATE/VERSION. 78/3/8.

1. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER. Rhintek, Inc.; Box 220; Columbia, MD 21045 (301).
2. MACHINE. Data General Nova or Eclipse minicomputers or equivalents. We are using the compiler on a Nova $3 / \mathrm{D}$ running Rev. 6.10 mapped RDOS. However, we are cleaning up the code and expect the compiler to be able to run under unmapped RDOS on a 32 k Nova within a few weeks.
3. SYSTEM CONFIGURATION. Mapped RDOS system or 32 k unmapped RDOS with minimum operating system. The current revision of Data General RDOS will be supported but the compiler should work with older versions
4. DISTRIbUTION. 9 track magnetic tape, $800 \mathrm{bpi}, 7.5$ inch tape in the RDOS dump format. Price for a single user license is $\$ 975$. Multi-use, OEnís, and educational use licenses will be handled on a separate basis.
5. DOCUMENTATION. The package includes source code, binary code, and ready-to-run demo programs. Instructions for executing the compiler are included; the operational
6. MAINTENANCE. Updates for 1 year and notification of substantial enhancements as long as interest is shown. We will maintain a users group and encourage bug reports and suggestions. This compiler is used by Rhintek as an application and system programing language and will continue io receive support and enhancements by us
7. STANDARD. Based on Sequential Pascal which varies from Standard Pascal. The current version lacks: file, goto, label, and packed reserved words and sqr, sin, cos, arctan, in, version lacks: file, goto, label, and packed reserved wo
exp, sqrt, eof, eoln, odd, and round built in functions.
. MEASUREMENTS. The compiler compiles source code at the rate of 200 lines/min. This is about $1 / 2$ the rate of the PDP-11/45 but about 5 to 10 times the speed of other compilers on the Nova. The compiler will compile itself in about 30 minutes total. (* Compilation and execution space requirements not reported. *)
8. RELIABILITY. Good. (* date first released, number of sites using system not reported. *)
9. DEVELOPMENT METHOD. The virtal machine was coded in Nova assembler language and then the compiler was modified along with the interpreter into its present form. (* Person-months to develop system not reported. *)
10. LIBRARY SUPPORT. There is no library support as yet. The operating programs support program swapping or chaining with only minimal effort as this is used with the compiler.

Data General Eclipse San Bernadino

MDC ECLIPSE RDOS PASCAL
Version 3
0. PRODUCT DESCRIPTION. MDC PASCAL Version 3 is an efficient PASCAL compiler and runtime support system designed for the execution of small PASCAL programs in a minicomputer environment. The development criteria are as follows:
A. To support interactive I/O in a reasonable way
B. To be compatible with, as far as possible, the existing MDC ECLIPSE RDOS PASCAL Version 2.
C. Close agreement with the P4 'standard'
D. A reasonable integration into RDOS. (We support background/foreground,
subdirectories, and a simple command-line form of activation)
E. Speed of execution is a primary concern in Version 3. The size of the object
program is secondary to this speed criterion.
the modularity and intelligibility of the code
The magnetic tape we distribute contains executable object code, source code, and machine readable documentation. It is assumed that the user has an existing MDC ECLIPSE RDOS PASCAL Version 2 operating at his site.

1. DISTRIBUTOR/IMPLEMENTOR/MAINTAINER. Ted C. Park; Director, Systems Development; Medical Data Consultants; 1894 Commercenter West, Suite 302; San Bernardino, CA 92408.
2. MACHINE. Data General - any ECLIPSE-Tine computer.
3. SYSTEM CONFIGURATION. ECLIPSE must have FPU or EAU, minimum of 16 K words user memory, RDOS REV 6.1 or greater, FORTRAN 5 (any recent revision).
4. DISTRIBUTION. System supplied on 9-track 800 BPI tape in RDOS 'dump' format. The cost is $\$ 100.00$ to cover our mailing and duplicating costs.
5. DOCUMENTATION. User must obtain his own copy of the Pascal Users Manual and Report. It is recommended that the user obtain an implementation kit from the University of Colorado. Documentation and operating procedures are supplied on the tape.
6. MAINTENANCE POLICY. Bug reports are welcome but no formal commitment for support can be made at this time. Extensive testing of the product has been done and all known bugs have been eliminated.
7. STANDARD. PASCAL P4 subset.
8. MEASUREMENTS.

Compilation Speed: Word Size:
Real Arithmetic:
Integer Arithmetic:
Set Size:
Minimum Memory Needed:
40 chars/sec (including blanks and comments)
16 bits
Uses 32 bit
Uses 16 bits
64 bits
Approximately the same as the code produced by Data General FORTRAN V compiler 16K words
9. RELIABILITY. Version 1 exists in at least 10 sites, we believe no bugs exits. Version 2 is primarily the same as Version 1 except with improved operating procedures faster compiles and executions, and increased capability; it also exists in at least 10 sites, we believe no bugs exist here either. Version 3 is a new product and has had thorough in-house testing. From our past experience, we have every reason to expect good performance in the field.
10. DEVELOPMENT METHOD. Developed from PASCAL-P4. Version 3 consists of a smal1 program which rearranges the PCODE output by the compiler into a form syntacticaliy acceptable to the Data General macro-assembler. A macro-library is supplied which output from the assembler may then be submitted to the normal Data General relocatin load procedure to produce an executable core image file. A runtime support library which includes some initialization routines, an error routine, I/O routines, and transcendental function routines is also included. All programs are written in assembly language and are extremely modular and well documented so that any changes wished by the user should be easy to incorporate
11. LIBRARY SUPPORT. No Data General libraries are needed to run the system nor is it possible to use any if desired.

Data General Nova Austin, TX
Department of Computer Sciences
Painter Hall 3.28
THE UNIVERSITY OF TEXAS AT AUSTIN
college of natural sciences
14 May 1978 AUSTIN, TEXAS 787 I2

Dear Andy,
I am enclosing three reports on work which $I$ have been doing (did) on implementing Pascal (or a Pascal-like. language at least) on a Nova 3/D This work differs from the University of Lancaster Version by directly compiling assembly code, not hypothetical stack code which must then be interpreted.
(* See Abstracts, above right. *)
Sincerely,

"Using Pascal on the Novas"
Abstract: This note describes the procedure for using the Pascal compiler on the Nova mputer system at the Department of Computer Science at the University of Texas at Austin. It also indicat es the limitations of the system and how they can be overcome.
"A Compiler for a Pascal-like Language"
Abstract: The development of major software systems for the Nova computer system can benefit greatly from the existance of a systems programming language. The development of such a language, and its supporting compiler is currently underway. This note reports on the language definition ant the mechanics of the compiler.

Abstract: A compiler is being written to translate a Pascal-like language into assembly code for the Data General Nova 3/D computer. A previous note has described the language and the basic structure of the compiler. In this note, we describe the code-generation problems encountered and their solution.

Data General Nova 840 Barcelona, Spain

We have received no new information on this implementation since that which we published last year in Pascal News issue: \#9-10: 81-82.

Data General Nova (Lancaster)
July 27, 1978
Dear Andy,
Enclosed is my renewal and here is some up-to-date information on our PASCAL distribution effort:

We are currently distributing Revision 2.01 of the Lancaster compiler for the NOVA. This revision has el iminated some of the minor problems found in the first release and has added some enhance-
ments, such as separately compiled procedures and support for random I/O. The source code and binaries are available on magnetic tape for $\$ 140$. The binaries on 1 y are $\$ 70$.

We have had a tremendous response to our press releases about the compiler, and have shipped some 59 copies so far, including copies sent to 7 foreign countries, even though we are only soliciting U.S. business. The number of reader response bingo-card inquiries is approaching 1000, indicating a high degree of interest in the language, many of the inquiries have come from England where this version was originally developed. Our customers have had very few problems with the Lancaster software, and we now have several applications programs running in PASCAL on the NOVA.

Sincerely,
Hank
H. S. Magnuski

Gamma Technology, Inc.

## GAMMA TECHNOLOGY


0. DATE/VERSION. Checklist last updated 77/10/27.

1. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER.
(Europe, Asia, Africa):
R. E. Berry and A. Foster;

Dept. of Computer Studies
University of Lancaster;
Bailrigg, Lancaster LAl 4YX, U.K.;
H. S. Magnuski
H. S. Magnuski

800 Welch Rd.
Palo Alto, CA 94304
TWX: 910-373-1296

Implementors: R. E. Berry and A. Foster.
MACHINE. Data General Nova series $(2 / 10,820)$.
3. SYSTEM CONFIGURATION. RDOS $4.02 / 5.00$ operating system; 32 K core, disk backing store. No hardware multiply/divide or floating point needed. One user reports using system with RDOS without any trouble.
4. DISTRIBUTION. From Lancaster: Cassette tape or 2.5 Mbyte cartridge disk (* cost not reported *) From Palo Alto: 800 bpi 9 track tape, binary only- $\$ 70$, with source- $\$ 140$; From Swampscott: (* format, cost not reported *).
5. DOCUMENTATION. A 82-page user manual is provided. (* Not known if this is machine retrievable. *)
6. MAINTENANCE. No formal commitment to provide support can be given, however, bug reports•are welcome. To date all known bugs have been fixed and this policy will continue as long as is practicable.
7. STANDARD. Pascal P4 subset accepted. Extensions for random I/O provided.
8. MEASUREMENTS. Typical runtimes compare favorably with those of other languages generally available on the Nova. P-code is generated, assembled and then interpreted
Compiler NMAX (decimal)
Release
elease 2
additional fixed table space 1,092 1,197 (in words)
The workspace remaining depends upon size of the RDOS system used. The size of program which can be compiled depends on the number of user defined symbols (dyunamic area used) and depth of nesting of procedures/statements. Thus it is difficult to make any general
statement about the size of program which can be compiled, however, we observe that the assembler for the system is some 1,100 lines of Pascal source generating 7,400 P-code instructions and we can compile this on our 32 k system. We cannot compile the compiler but would expect to do so with more than 32 k core

Timing information for Nova Pascal Lancaster Release 2: We have not yet compiled the compiler with our system so we cannot give figures for that. Instead to provide the basis for our statement that the performance of our Pascal "compares favorably" with DG ALGOL a
list of times obtained by running some well known small, and often uninteresting program are given. The timings are taken from a Nova $2 / 10$ running under RDOS 4.02 with 32 k of core an no hardware multiply/divide or floating point. They were (rather crudely) obtained by using the GTOD command to. prefix and postfix the CLI command necessary to load the appropriate program. "Compile" should be taken to mean the production of a save file (.SV) from the source program.
Programs:

1) A program consisting simply of begin end.
2) Matrix Mutiply of two $50 \times 50$ integer matrices (no I/O).
3) Matrix Nutiply of two $50 \times 50$ real matrices (no I/O)
( I )
4) Ackermans function ( 3,6 ) (no I/O).
5) Read 10,001 integers from a file.
6) Generate 5000 random integers (printing only the last).
7) Generate 5000 random integers and write to a file.

|  | ALGOL |  | Pascal |  |
| :--- | :---: | ---: | :---: | ---: |
|  | compile | run | compile | run |
| $\# 1$ | $: 55$ | $: 06$ | $1: 31$ | $: 07$ |
| $\# 2$ | $1: 15$ | $1: 54$ | $1: 39$ | $2: 35$ |
| $\# 3$ | $1: 16$ | $14: 32$ | $1: 40$ | $11: 59$ |
| $\# 4$ | $1: 10$ | $2: 06$ | $1: 38$ | $5: 56$ |
| $\# 5$ | $1: 09$ | $2: 52$ | $1: 37$ | $1: 55$ |
| $\# 6$ | $1: 06$ | $3: 18$ | $1: 35$ | $1: 11$ |
| $\# \# 7$ | $1: 08$ | $1: 28$ | $1: 36$ | $1: 03$ |
| $\# 8$ | $1: 36$ | $1: 56$ | $1: 57$ | $3: 13$ |
| $\# \# 9$ | $1: 36$ | $4: 46$ | $1: 57$ | $4: 30$ |

Timings such as these offer much scope for debate. It is safer to let others draw what conclusions they will from these figures (and from any other figures which may be produced). I simply wish to observe that interpretive Pascal "compares favorably" with the code produced by DG ALGOL. In the programs used above the ALGOL and the Pascal look very
nuch the same. No attempt is made to exploit one feature of a particular language or mplementation, and no tuning has been done. If anyone has other examples to contributr to such timing comparisons I would be glad to hear about them.
9. RELIABILITY. Release 2.01 has been distributed to 50 known sites. No significant bugs have been reported from external users. First released 77/01; Latest release 78/7/27.
10. DEVELOPMENT METHOD. Originally cross-compiled from a CDC 7600. The P-code assembler was written from scratch in Pascal; the P-code interpreter was implemented in Nova was written from scratch in Pascal; the P-code interpreter was
assembly language. ( $*$ Person-months to create system not reported. *)
11. LIBRARY SUPPORT. No library support in release 1. Under Release 2 user procedures may be separately compiled enabling
to link into any other libraries.

DEC -- Introduction

## University of Montana <br> Missoula, Montana 59812

DEPARTMENT OF COMPUTER SCIENCE
Phone: (406) 243-2883

Dear Andy:
The DECUS PASCAL SIG is alive and well even though I am now in the Big Sky Country (Montana). My steering committee now resides in the four corners of the United States, but we are actively working on several PASCAL related
projects. We are keeping in touch with Seved Torstendahl (Sweden) as a US focal point for his PDP-11 PASCAL Compiler. In addition, we are actively pursuing the implementation of the NBS (National Bureau of Standards) PASCAL Compiler on the following PDP-11 operating systems: UNIX, RSX-11, IAS, RSTS, and RT-11. In addition to PDP-11's a smal1 portion of our group is working on a version of the NBS PASCAL Compiler for the VAX-11/780. We are very inpart of Ken Bowles' meeting at UCSD this summer and Justin Walker (NBS) is interested in implementing some of the agreed upon extensions for externall compiled modules. Please publish as much of the UCSD summer meeting report as possible in future issues of the PUG newsletter.

Dr. Roy Touzeau, also of the Computer Science Department here at the University of Montana, is al SO Working on a DECSYSTEM-20 version of Charles Hedrick's DEC-10 (KLIO) PASCAL Compiler from Rutgers University. He has mod-
ified the run-time system to remove the dynamic page management code as the ified the run-time system to remove the dynamic page management code as the
DEC-20 does its own paging. He is presently changing the run-time support to DEC-20 does its own paging. He is presently changing the run-time support to code. Future plans are to produce a one-step compiler/linker for student use in introductory programming courses. Any comments or suggestions regarding this effort may be sent directly to Roy.

Sincerely yours,


DEC LSI-11 UCSD
John R. Barr


We have received copies of two papers on the UCSD Pascal system; the titles are: "A Brief Description of the UCSD Pascal Software System' (*78/6/1*), and "Newsletter \#2--UCSD Pascal Project" ( $78 / 5 / 30$ *).

Jim McCord; 330 Vereda Leyenda; Goleta, CA 93017; 805/968-6681 reports: "I am acting as the distributor for UCSD Pascal for hobby users of the LSI-11. Cost is $\$ 50$, of which \$35 goes to UCSD for continued work. Other $\$ 15$ pays for documentaion and postage, if user
sends me 4 floppies. (Else I will provide for $\$ 3$ each.) This includes all source code for everything, including the interpreter. Anybody interested should get in touch with me (we already have 7 users).

Following checklist submitted by George Gonzalez, Special Interactive Computing Lahoratory; 134 Space Science Center; Jniversity of Minnesota; Minneapolis; MN 55455 on 78/10/01.
0. DATE/VERSION. I.4, released about May, 1978.

1. IMPLEMENIOR/DISTRIBUTOR/MAINTAINER. UCSD Pascal Project; Institute for Information Studies; University of California-San Diego; Mail Code C-021; La Jolla, CA 92093; 715/452-4526.
2. MACHINE. PDP-11, LSI-11 series with $16-28$ kwords memory; and various 8 and 16 -bit micros: Intel 8080, Zilog Z-80, etc.
3. SYSTEM CONFIGURATION. Has own operating system. Does not run under any other system (but can be brought up under $\mathrm{CP} / \mathrm{M}$ ). Requires $16-28$ kwords (unmapped).
. DISTRIBUTION. Source \& object programs available on RXOl diskettes. Contact UCSD for more information. Cost - $\$ 50$ for binaries; $\$ 200$ for source, maintenance and binaries.
4. DOCUMENTATION. User Manual. Gives overview of operating system and differences ith/extensions to Standard Pascal. Not machine retrievable.
5. MAINTENANCE. One-year maintenance (optional at higher cost)
6. STANDARD. Not implemented: Program header with file parameters; procedures dispose, pack, unpack; no procedures or functions as parameters; no boolean conversion in write procedure;

Differences: input is initially undefined; read(input,ch) is defined as begin get(input); ch := input ${ }^{\sim}$ end, instead of the Standard Pascal definition; rewrite requires a second parameter which specifies the system file name; files are not automatically osed at block exit; gotos cannot cross block boundries

Extensions: Numerous (but ill-defined) extensions: character strings as an intrinsic ype; string-manipulation facilities; random access to files; dynamic file opening/closing; shared variables for system communication; I/O error detection
capability; segmentation (overlay) scheme.
8. MEASUREMENTS. Compiles a $3400-1$ ine program in 28 k words, at $400-600$ lines /minute. (* How this compares with FORTRAN, other languages not reported. *) (* Execution speed, space not reported. $\star$ )
9. RELIABILITY.

The reliability of the Standard Pascal constructs is good.
Large ( 3000 line) programs, plus several 'portable' Pascal programs (XREF, COMPARE, PRETTYPRINT) have been run with no problems attributable to the Standard Pascal constructs.

The reliability of the UCSD "extensions" is generally poor.
The string-manipulation intrinsics (COPY, POS, CONCAT) do insufficient error checking. The graphics intrinsics do not check for out-of-range arguments (which usually crash the program). Writing on a reset'ed file can destroy other files. The compiler change the value of the literal. Writing a file which overflows awailable space does not cause an error.
10. DEVELOPMENT METHOD. P-code compiler/interpreter system. Based on P2. First released 77/8/1. About 300 sites using system.
11. LIBRARY SUPPORT. Compiler can read external source files. Predefined procedures ar provided for text-string manipulation, memory-mapped graphics, and system level input/output. These intrinsics are generally ill-defined and unreliable. No symbolic dump is available. The object-code level debugger supplied requires extensive knowledge of th

DEC PDP-8 (Minnesota)
O. DATE/VERSION. Checklist updated 78/10/5.

1. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER. John T. Easton, 612/373-7525; James F. Miner 612/373-9916; Address correspondence to: Pascal Group; SSRFC; 25 Blegen Hall; University of Minnesota; 269 19th Ave. South; Minneapolis, MN 55455; 612/373-5599.
2. MACHINE. Digital Equipment Corp. PDP-8/e.
3. SYSTEM CONFIGURATION.

OS/8 version 3. Hardware required:
-12 K minimum of core/RAM. 32 K is required for compilation. Can use up to 128 K
4. DISTRIBUTION. Release scheduled for second quarter, 1979
5. DOCUMENTATION. Machine-retrievable supplement to Pascal User Manual and Report (about 25 pages), in preparation.
6. MAINTENANCE. A policy has not yet been determined.
7. STANDARD. Emphasis has been on close adherance to the Pascal User Manual and Report There are two major restrictions: a) Procedures and functions may not be passed a parameters. This restriction will not be lifted without full type checking (which requires a change in the Pascal Standard). b) Files may be declared only in the main program, and files may not be components of arrays, records, or files; nor may files be allocated with the procedure NEW. Minor restrictions: set size $=96$ elements; maxint $=8,388,607 \quad$ ( $2 * * 23-1$ ) b) default
8. MEASUREMENTS

Execution speed--roughly comparable to FORTRAN IV (F4). I/O tends to be faster than Execution space--Interpreter takes 8 K , space needed for P -code and runtime storage depends on program.

9．RELIABILITY．Fair to good and improving．An earlier implementation has been in use a
10．DEVELOPMENT METHOD．As with most languages on the PDP－8，Pascal makes use of an interpreter（a modification of P－code）written in MACREL．The compiler（about 5000 lines， based on Pascal－P4）is written in Pascal．All standard procedures are written in MACREL． The implementation is not suitable for real－time applications．

11．LIBRARY SUPPORT．Currently（78／11／15），none planned for the first release．

DEC PDP－11（Amsterdam）

0．DATE／VERSION．Checklist not updated since 78／02．
1．TMPLEMENTOR／DISTRIBUTOR／MAINTAINER．Sources，binaries，and documentation are part of the third UNIX software distribution．Implementor：Johan Stevenson，Vrije Universiteit Maintainer：Andrew S．Tanenbaum；Vakgroep Informatica；Wiskundig Seminarium，Vrije Universiteit；De Boelelaan 1081；Amsterdam，The Netherlands；020／548－2410．
2．MACHINE．Any PDP－11 on which UNIX version 6 will run．
3．System configuration．See 2.
4．DISTRIBUTION．Through the UNIX software distribution center．，（＊No information on cost reported．＊）
5．DOCUMENTATION．Short manuals for the compiler and interprater in UNIX MAN format and a 12 page description giving details about the implementation．
6．MAINTENANCE．Bug reports are welcome．There will be an improved release of the current system．However，we are working on a totally new one．Main differences from the old one are：
anew hypothetical stack computer named EMI（see Tanenbaum，A．S．，＂Implications
of structured programming for machine architecture＂，CACM，Dec．1977）．
or lntermediate machine allows very compact code（only 15，000 8－bit bytes
microprogrammable computer must be easy．Moreover，this EM1 machine
microprogrammable computer must be easy．Moreover，this
allows compilation of other high level languages as well．
－an new interpreter with all kinds of．run－time checks and debugging aids．
expansion on EMI codes into PDP－11 instructions．

7．STANDARD．Main differances with Standard Pascal are：
－no gotos out of procedures and functions．
procedures and functions can not be passed as parameters．
－extern procedures and functions not implemented．
－mark and release instead of dispose．
－at most 8 files（all text），including input and output．
－An explicit get or readn is needed to initialize the file window
procedure unpack not availlable，packed fign in record declarations
procedure unpack not available，packed ignored；all records are automatically packed． between 0 and 63 inclusive．maxint $=32,767 .(2 * * 15-1)$ ．Setsize $=0 . .63$ ．Full ASCII accepted（parity ignored）．Keywords and standard names are recognized in lower case．

## 8．MEASUREMENTS．

compilation speed－$-40,000 \mathrm{char} / \mathrm{min}$ on a $11 / 45$ with cache
compilation space－48k bytes to compile the compiler．Very big programs can be compiled． execution speed－－you lose a factor of 8 by interpretation．However，I／0
is relatively fast．Compared to interpreted Pascal on
a big machine（CDC Cyber 73）it is 10 times slower．
The binary code of the compiler is 23,000 bytes．

9．RELIABILITY．The compiler and interpreter are good．However，the run－time checking of the interpreter is poor．Preliminary version first ran in 1977．（＊Date system first released to users，number of sites using system not reported．＊）
10．DEVELOPMENT METHOD．The compiler is based on the Pascal－P2 compiler．A Cyber 73 was used for bootstrapping．The time needed by one inexperienced implementor was about 6 months．

11．LIBRARY SUPPORT．No library support at all．There are some hidden library routines used by the system．

DEC PDP－11 Berkeley

UNIVERSITY OF CALIFORNIA，BERKELEY


PROGRAM IN OUANTTTATIVE ANTIROPOLOGY
DEPATTMENT OF ANTHROLOLOGY

2290 PIEDMONT AVENUE ${ }^{2}$
29 April 78

Dear Andy，
I was suprised that there wann＇t anything in the PN last time about the Berkeley UNIX（PDP－11）Pascal．I thought I＇d let you know it exists，since the implementors apparently haven told you anything．

It is an interpretive system written for support of computer science instruction，so it is very fast at generating（intermediate）code，but slow at execution．The syntax scan is the best I＇ve seen（of any compiler for any language） it is very informative for unexperienced users，comments on suspicious（but syntactly ELSE．Such corrections show on the listing but the correct intermediate code is generated－the note will continue to appear on subsequent listings until the source file is changed by the user，of course．Definately accepts Standard Pascal： I swap very large programs back and forth between the PDP－11 and the CDC 6400 with only changes required in first and last character constants（MINCHAR and MAXCHAR）．

The development was supported at least in part by US ERDA，and the authors seem willing to distribute it for instructional use．A fifty－one page user＇s manual， the Computer Science Library for a couple of bucks．The authors of the manual are William N．Joy，Susan L．Graham and Charles B．Haley．Joy and Graham can be reached at the UCB Computer Science Division，Department of Electrical Engineerin and Computer Science，University of California，Berkeley，Berkeley，CA 94720. Graham＇s office phone 非is 4I5－642－2059．I think Haley has left，I have a vague recollection that he is at Bell Labs now．

This is an exeellent Pascal system，which I would recommend highly to anyone running under UNIX．Of course，since it is an interpretive system ther
would be execution time problems for some production applications．
p．s．Runs on $11 / 45$ and 11／70．Doesn＇t accept procedure and function names as parameters．I＇ll send you some documentation if I get time．


We have received no new information on this UNIX, RT-11, DOS, and RSX-11 We have received no new information on this UNIX, RT-11, DOS, and RSX
implementation since that which we published last year in Pascal News issue: \#9-10: 83 . DEC PDP-11 Missolla, MT

We have received no new information on this RSX-11 implementation since that which we published last year in Pascal News issue: \#11: 91.

DEC PDP-11 (OMSI) (formerly ESI)
Maurice R. Munsie; Network Computer Services P/I; 69 Clarence St.; Sydney 2000 Australia reports: "We are distributing in Australia OMSI Pascal-1. A number of sales have
already been made and plans are being made for the OMSI implementors to hold workshops in Australia later this year." (* 78/8/28*)
0. DATE/VERSION. 77/12/76.

1. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER. Oregon Minicomputer Software, Inc. (OMSI); 4015 SW Canyon Road; Portland, OR 97221 ; 503/226-7760. Implementors: John Ankcorn, Don Baccus, and Dave Rowlar .
2. MACHINE. Any model Digital Equipment Corp. PDP-11 or LSI-11.
3. SYSTEM CONFIGURATION. Minimum of 16 K words. Operates under RT-11, RSTS/E, or RSX.
4. DISTRIBUTION. Compiler, support module, cross referencer, text editor and instruction manual available for $\$ 1500$ ( $\$ 995$ for educational use). Available on 9 track 800 bpi magnetic tape, or DEC cartridge disk.
5. DOCUMENTATION. Over 70 -page machine-retrievable instruction manual. Currently 5. DOCUMENTATION. Over ${ }^{70}$.
(76/11/02) working on more.
6. MAINTENANCE. One year of unlimited fixes and updates, followed by annual subscription service. (* Reported by users that "vendor seems to be responsive in terms of support". *)
7. STANDARD. Full standard plus extensions: additional features for real-time hardware control; separate compilation of procedures; Macro (assembler) code in-line insertion; actual core addresses of variables can be fixed (giving access to external page I/O addresses at the Pascal level.
8. MEASUREMENTS
compilation speed-About 3500 characters /second, on the PDP-11 model 05. compilation space--very economical-it can compile 3000 line programs 28 K on PDP-11/40. No overlays are used in the system. bout twice as fast as the DEC FORTRAN IV and many times faster than DEC BASIC. A worst-case 'number-cruncher' example ran at $40 \%$ faster than the DEC original FORTRAN.
execution space--very economical-much of the space improvement over DEC
FORTRAN is due to the smaller support module for Pascal.
9. RELIABILITY. Excellent--far better than DEC FORTRAN. In use since 75/11. Over 100 installations, and growing steadily.
10. DEVELOPMENT METHOD. Single-pass recursive-descent compiler written in Macro-11. Hand-coded based on University of Illinois bootstrap (with extensive changes) in about two person-years of effort. First compiler written by both implementors. Compiler translates source into Macro-11 which is then assembled and linked to the support module for execution.
11. LIBRARY SUPPORT. Separate compilation of procedures with load-time insertion and linkage is implemented.

DEC PDP-11 Redondo Beach

We have received no new information on this Concurrent Pascal (SOLO) implementation since that which we published last year in Pascal News issues: \#11: 89-90.

DEC PDP-11 (Stockholm)
0. DATE/VERSION. 77/12/22.

1. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER. Seved Torstendah1; Tn/X/Tdg.; Telefon $A B$ LM Ericsson; AL/Ufe; S-125 26 Stockholm, Sweden; 08/719-4909.
2. MACHINE. Digital Equipment Corp.:

DEC-10 (cross-compiler that generates code for all PDP-11's);
PDP-11 model 35 and up (self compiles
and generates code for all PDP-11's);
he compilers generate code for floating point hardware and extended arithemtic instruction sets if option switches are set.
3. SYSTEM CONFIGURATION. DEC-10 cross-compiler: TOPS-10. PDP-11: RSX-11M (Probably it is an easy task to replace the RSX interfacing routines with new ones interfacing to DOS or
RT-11; but we do not plan to do that work here. Maybe routines to interface with RSX-11S RT-11; but we do not plan to do that work here. Maybe routines to interface with RSX-11S
will be made.) PDP-11 with memory management and a user partition of at least 28 k words, will be made.) PDP-11 with memory management and a user partition of at least 28k words, preferably 32 k words.
4. DISTRIBUTION. The compilers are available at $\$ 50$, plus $\$ 10$ if we supply the tape ( 600 feet). The distribution set includes source and object modules of the compilers and the runtime library, command files following formats; indicate which you want:

- three DECtapes in PDP-11 DOS format (DEC-10 and PDP-11 users)
- one 9-track magnetic tape in DEC-10 format (DEC-10 users)
- one 9-track magnetic tape in industry standard format
- one 9-track magnetic tape in DOS format (PDP-11 users).

5. DOCUMENTATION. A machine-retrievable user manual, complementing the Pascal User Manual and Report, is included on the distribution tape.
6. MAINTENANCE. No responsibility, but if errors are found reports will be distributed to known users. Error reports and improvement suggestions accepted.
7. STANDARD. With regard to the definition of Pascal in Pascal User Manual and Report, the following restrictions hold:

- packed data structures are only implemented for character arrays (always packed, two chars/word) and for Boolean arrays (packing optional, one Boolean / bit). The standard procedures pack and unpack are not implemented.
- only local jumps are allowed
- a pair of procedures, "mark" and "release", have been added to allocate and deallocate dynamic storage.
The following extensions have been implemented:
- function results can be of a nonscalar type.
- arrays with unspecified bounds (but specified index-structure) can be used as formal parameters to. procedures, allowing differently declared variables or constants to be
a string parameter type h
or substrings thereof may be been introduced in which one-dimensional character arrays ent characters are considered as "read-only".
- procedures may be compiled separately. the procedure block replaced with "extern".
- most option selectors ( ( $* \$ M+*$ ), etc.) are selectable by switches on the MCR command line (version 5, 77/12).

8. MEASUREMENTS.
compilation speed--about 300 characters/second; increases to 3000 characters/second in a 64 k words partion using
PLAS under RSX-11M.
compilation space--The compiler requires a 32 k word partion (at least 26 k words for very small programs).
execution speed--(* No information provided. *)
execution space--(* No information provided. *
(* How this compares to FORTRAN and other languages not reported. *)
9. RELIABILITY. Excellent. The compiler is now in use at over 200 sites. Only minor errors have been found since July, 1977. First version released April, 1977. Latest version: December, 1977.
10. DEVELOPMENT METHOD. The compiler is a modification of the cross compiler from Mr. Bron, et. al. of Twente University of Technology in the Netherlands. The original cross-compiler was written in Pascal and developed from Pascal-P. Two major modifications have been undertaken:
the compiler generates standard object modules

- the compiler gives full access to the RSX/IAS file system.

The compilers are written in Pascal, and both have the same source code except for two
separately compiled routines. The cross compiler is generated when the DEC-10 pascal compiler from Hamburg compiles the source. When it then compiles itself the PDP-1l version is created. The cross compiler for PDP-11 running on DEC-10 produced by Bron et al was used as input. This compiler was modified to generate object code linkable under RSX-11M and to give access to the file system of RSX-11M. When the cross compiler was finished it compiled itself and was thus transfered to the PDP-11. The implementation effort until now (77/02/09) has been about five person-months. To make use of floating point hardware another two person-months will be needed. A new version which performs some optimization will probably be developed later.
11. LIBRARY SUPPORT. Separate compilation allowed. Possible to use external procedures written in FORTRAN or assembler. The December 1977 version also gives: Automatic copy of
text from library into source program (include); execution frequency measurements; execution trace; option selectors ( $(* \$ R-*)$, etc.), settable by switches in the MCR command line. Next version (Spring, 1978) will also include a symbolic post-mortem dump an an interactive source-level debugging package (mainly copied from the DEC-10 Hamburg-DECUS compiler).

DEC PDP-11 Tampere, Finland

The DEC PDP-11 Stockholm Pascal system (for RSX-11M) was modified slightiy during October, 1977 to run under IAS by: Jyrki Tuomi and Matti Karinen; Tampere University of Technology; Computing Center; SF-33100 Tampere 10; Finland; (* No phone number reported *). A 60 -page report on this implementation (in Finnish) is available from Tampere.

## DEC PDP-11 Twente

We have received no new information on this implementation of a cross-compiler from DEC-10 to any PDP-11 on any operating system since that which we published last year in Pascal News issue: \#9-10: 85 .

We have received no new information on this RSX-11D implementation since that which we published last year in Pascal News issue: \#9-10: 85-86.

DEC VAX-11/780 Seattle

We have received no new information on this implementation since that which we published last year in Pascal News issue: $\# 12: 63$.

DEC VAX-11/780 (Redondo Beach)
We have heard rumors that an implementation is underway at TRW corporation at Redond Beach, CA.

## DEC-10 (Hamburg-DECUS)

O. DATE/VERSION. Checklist not updated since 77/08

1. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER. Implementor/Maintainer: E. Kisicki; H. -H. Nagel; Universtat Hamburg; Institut fur Informatik; SchluterstraBe 66-72; D-2000 Hamburg 13 Germany; 040-4123-41.51; TELEX: 214732 uni hh d.

Maynard, MA 01754
USA;
617/ 897-5111;
TWX: 7103470212
(Eastern Hemisphere):
DECUS-Europe;
P. O. Box 340;

CH-1211;
Geneva 26, Switzerland;
O22/ 4279 50;
2. MACHINE. Digital Equipment Corp. DEC-10. (Adapted to the DEC-20 by DEC).
3. SYSTEM CONFIGURATION. DEC TOPS- 10 moniter using Concise Command Language (CCL). U.ses KA-10 instruction set. Modifications to use KI-10 improved instruction set have been made by Charles Hedrick.
4. DISTRIBUTION. From DECUS (Digital Equipment Corp. User's Society).
5. DOCUMENTATION. Machine-retrievable manual included on distribution tape.
6. MALNTENANCE. No regular maintainance can be given.
7. STANDARD. Extensions: Functions FIRST and LAST for scalars; UPPERBOUND and LOWERBOUND for arrays; MIN and MAX available as standard functions; procedures to determine the value of CCL options available; otherwise in case statement; LOOP...EXIT IF...END statement; Initialization procedure.
8. MEASUREMENTS. (* No information provided. *)
9. RELIABILITY. Very good. First version released in 75/7. Distributed to at least 60 sites. Later version operational in 76/9. Latest version released to DECUS in 77/2.
10. DEVELOPMENT METHOD. Pascal-P2 and subsequent self bootstraps. Latest version dated 76/12/30.

11．LIBRARY SUPPORT．Symbolic post－mortem dump available．Interactive run－time ource－level debugging package available．Separate compilation and inclusion in relocatable object code library of Pascal，FORTRAN，COBOL \＆ALGOL，and MACRO－10 assembler routines．

## DEC－10 Systems－Pascal

We have received no new information on this implementation since that which we published last year in Pascal News issue：\＃9－10：90－91．

DEC－20

See DEC－10 Hamburg－DECUS

## Dietz Mincal 621 Hamburg

We have received no new information on this implementation since that which we published last year in Pascal News issue：$⿰ ⿰ 三 丨 ⿰ 丨 三 ⿻ ⿻ 一 𠃋 十 一 八$－10：91－92．

## FOXBORO Fox－1

We have received no new information on this implementation since that which we published last year in Pascal News issue：\＃9－10： 92.

## fuIITSU Facom 230－30 Tokyo

We have received no new information on this implementation since that which we published last year in Pascal News issue：\＃9－10： 92

## FUJITSU Facom 230－55

See FUJITSU Facom 230－30 Tokyo．

## Harris／4 Delft

We have received no new information on this implementation since that which we published last year in Pascal News issue：\＃9－10： 92.

## Heathkit H－11

（＊This machine is based on the LSI－11 microprocessor from DEC and it is believed that the DEC LSI－11（UCSD）implementation will run on this machine；though nothing
definite has been reported．＊）
According to Bill Schiffbauer；Sales Coordinator，Computer Products；Heath Company Benton Harbor，MI 49022；616／982－3285；TELEX 72－9421：＂At this time（＊77／11／15＊），Heat has no plans to offer a Pascal compiler or interpreter．．．Since the H－11 uses the LSI－11， the［UCSD Pascal］compiler should be compatible with the H－11．

According to Robert W．Furtaw；Marketing，Heath Company，Benton Harbor，MI 49022：（＊ $78 / 1 / 19 *$ ）＂We also have been observing the appeals for Pascal appearing in recent publications．However，we presently have no immediate plans to offer one for our system．
With all the interest，I would not be sufprised to see one which could easily b reassembled for our system．＂

## Hewlett Packard HP－2100（Trieste，Italy）

0．DATE／VERSION．78／10／9．
1．IMPLEMENTOR／DISTRTBUTOR／MAINTAINER．Implementor：Paolo Sipala；Instituto di Electrotechnica；Universita di Trieste；Via Valerio，10，34127；Trieste，Italy Tel．040－733033．Distributor：Hewlett－Packard Software Center；Contributors Section； 11000 Wolfe Blvd．；Cupertino，CA 95014；（＊No phone number reported．＊）
2．MACHINE．Hewlett Packard HP－2100 or 21MX
3．SYSTEM CONFIGURATION．Old version－DOS IIIb；New version－RTE．There are seperate versions for EAU，non－EAU，and floating point hardware．Requires an llk main area．

4．DISTRIBUTION．（＊No information reported on cost，distribution formats．＊）
5．DOCUMENTATION．（＊No information provided．＊）
6．MAINTENANCE．（＊No information provided．＊）
7．StANDARD．（＊No information provided．＊）
8．MEASUREMENTS．Requires an 11 k main core area（so it might fit in a 16 k system，if the resident operating system modules are kept to a minimum，but 24 k is more comfortable）．It is not noticably slower than the standard compilers when compiling，and not worse than the standard interpreter（BASIC）when interpreting．

9．RELIABILITY．Has been subjected to rather limited testing（a few dozen programs fro the Users Manuai）and is now（＊78／3／20＊）being offered to students for their use．
10．DEVELOPMENT METHOD．A P－code interpreter written in HP－Algol．
11．LIBRARY SUPPORT．（＊No information provided．＊）

## Hewlett Packard HP－21 MX Durban

We have received no new information on this implementation since that which we published last year in Pascal News issue：\＃9－10： 93 ．

Hewlett Packard 3000 Santa Clara

We have received no new information on this implementation since that which we published last year in Pascal News issue：\＃9－10： 94 ．

## Hewlett Packard 3000 Sunnyvale

We have received no new information on this implementation since that which we published last year in Pascal News issue: \#12: 63-64.

## HITACHI Hitac 8800/8700 Tokyo

See also IBM 360/370. We have received no new information on this implementation since that which we published last year in Pascal News issue: \#9-10: 94.

Honeywell 6000, level 66 (Water1oo)
0. DATE/VERSION. Checklist not updated since 77/08.

1. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER. Implementor: W. Morven Gentleman; Mathematics Faculty Computing Facilty; University of Waterloo; Waterloo, ONT. N2L 3G1; CANADA; 519/ 885-1211. Distributor: Honeywell Information Systems; Waltham, MA (* See local HIS sales
office first. *)
2. MACHINE. Honeywell 6000 series; level 60/66. Operates under GCOS (TSS). Currently (* 76/03/08 *) a DRL TASK version is under consideration.
3. SYSTEM CONFIGURATION. Honeywell level 66 or 6000 series with EIS. Minimum of 26 k words.
4. DISTRIBUTION. (* No information provided. *) (* Rumor has it that distributor charges extra for maintenance. *) 5. DOCUMENTATION. From Honeywell Information Systems; Publication Dept.; MS-339; 40 Guest
St.; Brighton, MA 02135: "A Pascal Product Brief", (\#AW66, free), 2 pg. (marketing oriented) and "Pascal User's Guide", (\#AW65, \$1.30), 30 pg . (reference manual). Machine retrievable supplement to Pascal' User Manual and Report; also includes extensions, restrictions, known bugs, etc.--about 45 pages total.
5. MAINTENANCE. Supported by University of Waterloo through agreement with HIS; some users have reported problems in getting Honeywell to pass bug reports on to Waterloo. Extensions planned to allow extern to be GMAP, COBOL, ALGOL, PL/I, B, C, etc.
6. STANDARD. Restrictions
-Program statement not accepted, replaced by required procedure 'main'.
-No files with components of type file.
-Only files of type char may be read or written (with the standard
xtensions:
-Files may be opened dynamically.
-Extended file handing is available.
-Various procedures and functions al and FORTRAN procedures may be used.

- Optional left-toright evaluation provide access to operating system.

Oplonal lef boolean expressions and if statements
-Alternate Interactive I/O package available.
-Full upper/lower case capability.
8. MEASUREMENTS.
compilation space--minimum of 26 k words. Typical programs require less than 30 k words. ompilation speed--(* No information provided.
ds depending on the program and the
Pascal support routines required
(* No information provided. *)

* How this compares to FORTRAN and other languages not reported. *)
. RELIABILITY. (* No information provided on number of sites using system.*) Some users have reported problems with compiler reliability and responsiveness of distributor. See Pascal News \#11: 34-36, 92-93. Distributed since 76/05. Version 6 expected in 77/12.

10. DEVELOPMENT METHOD. Independant implementation (unrelated to Pascal-P or CDC 6000 Zurich compilers); written in "B", an implementation language and successor of BCPL.
11. LIBRARY SUPPORT. Separately compiled Pascal and FORTRAN routines may be saved and called from user specified libraries at run time. A post-mortem debugger is planned, but presently (* 76/10/25 *) far from being implemented.

Honeywell H316 Minnesota
0. DATE/VERSION. 78/7/4.
. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER. Robert A. Stryk; Honeywell Corp. Computer Science Center; 10701 Lyndale Ave. S.; Bloomington, MN 55424; 612/ 887-4356.
2. MACHINE. Honeywell H-316.
3. SYSTEM CONFIGURATION. 32 k , dual cartridge disks, line printer, 7 -track magnetic tape.
. DISTRIBUTION. 7-track tape with programs to bootstṛap from BOS 210. (* cost not reported. *)
5. DOCUMENTATION. Informal comments on 316 kernal implem entation.
6. MAINTENANCE. No known errors, no work planned. Bob reported on 78/7/4: "changing jobs--Distribution of H316 Concurrent Pascal very cloudy".
. STANDARD. A modified implementation of Concurrent Pascal, which varies from Standard ascal.
. MEASUREMENTS. SOLO system needs minimum of 40 k to execute compilers
9. RELIABILITY. No known errors. (* Date first released, number of sites using system not reported. *)
10. DEVELOPMENT METHOD. The H316 kernal imitates the PDP-11 reversed byte addressing which makes it compatible with the distribution tape but a bit slow in execution. The makes it compatible with $\operatorname{BOS} 210$. The kernal is written in DAP700.
11. LIBRARY SUPPORT. That provided by the SOLO system.

## IBM 1130

We have received no new information on this implementation since that which we published last year in Pascal News issue: \#9-10: 101 .

IBM 360/370 AAEC

We have received a copy of a report titled "Implementation of Pascal 8000 on IBM 360 and 370 Computers" (* $78 / 8 / 4$ *) which is available from the distributor.
0. Date 78/09/12

1. Implementors:
T. Hikita and K. Ishihata,

Dept. of Information Science,
University of Tokyo,
2-11-16 Yayoi
Bunkyo-ku токчo,
113 JAPAN.
(HITAC - 8000 Version)
G.W. Cox and J.M. Tobias
yystems Design Section,
AAEC Research Establishment,
SUTHERLAND, 2232,
N.S.W. AUSTRALIA
(IBM 360/370 Version)

Distributors/Maintainers:
G.W. Cox and J.M.Tobias
address as above
2. Machines:

IBM360 and IBM370 - compatible machines
3. System Configuration:

The compiler runs under any of the oS family of operating systems - i.e. MVT, MFT, VS1, VS2, SVS and MVS. A CMS interface is currently being developed, soon to be available. A minimal program can be compiled in 128K; the compiler requires about 220 K to compile itself.
4. Distribution:

Write to G.W. Cox and J.M. Tobias at AAEC to receive an ordex form. The cost is \$Aloo; there is no agreement to be signed. Two systems are supplied: a "compile-and-go" system which has its own compiled-code format, and a "linkage-editor" system which produces IBM-standard object modules. Both source and load modules for these systems are supplied - the compilers are written in pascal and the runtime support in 360 Assembler.

An implementation guide, plus machine-readable implementation JCL, and machine-readable documentation are also supplied:

The system is distributed on a new 600 ft . magnetic tape at a density of 800 or 1600 bpi ; the tape is supplied by the distrjbutor.
5. Documentation

Machine-readable documentation is in the form of a report comprising a summary of extensions to Standard Pascal plus a complete specification of the language as implemented.
6. Maintenance Policy.

No guarantee on maintenance is given; however we are anxious to receive bug reports and suggestions, and will do our best to fix any problems which may occur.
7. Standard.

The full standard is supported with finiteness in a few areas:

- maximum static procedure nesting depth is 6 .
- maximum set size is 64. (this precludes set of char.) It is hoped to increase this very soon.
- maximum number of procedures in a program is 256
- maximum size of compiled code in any one procedure depends on its static level: the main program may be up to 24 K , and this is reduced by 4 K for each increment of stàtic nesting level.
Significant extensions to the standard are in the following areas:
- Constant definitions for structured types. It is therefore possible
to have arrays, records and sets as constants
- A 'value' statement for variable initialisation
- A 'forall' statement of the form:
forall <control variable> in <expression> do <statement>
where <expression> is of type set.
- A 'loop' statement, specifying that a group of statements should be repeatedly executed until an 'event' is encountered. Control may then be transferred to a statement labelled by that event.
- The types of parameters of procedures or functions passed as parameters must be specified explicitly, and this enables the compiler to guarantee integrity.
- The 'type identifier', restriction in a procedure skeleton has been relaxed to allow 'type'.
- Functions 'pack' and 'unpack' are supported, as are packed structures in general.
- Exponentiation is fully supported, and is used via the double character symbol '**'.
- A 'type-change' function has been introduced that extends the role of 'chr' and 'ord'.
- Case-tag lists may now range over a number of constants, without explicitly having to list each oonstant.
The range is denoted by
<constant> .. <constant> Thus,

4,6..10,15,30..45
is now a valid case tag list
A default exit is also supplied which can be used-if none of the other tags match.
Other interesting features of the system are

- Procedure 'new' is fully supported, obtaining the minimum heap requirements as specified by variant tags. Procedures 'mark' and 'release' are also supported.
- Files may be external or local. Thus, structures such as 'array of files' are available. External files are named in the program statement, local files are not. Both external and local files may be declared in a procedure at any level.
- Text-files with RECFM of $F[B][S][A], V[B][S][A]$ and $U[A]$ are supported. Non-text files must have $\mathrm{RECFM}=\mathrm{F}[\mathrm{B}]$.
- All real arithmetic is in double precision ( 64 bit floating-point format).
- Control of input and output formatting is as described in the Jensen and Wirth report. The form is
$\mathrm{x}[\mathrm{n}] \quad[\mathrm{m}]$, where n and m are integer expressions.
Further, elements of type packed array of char may be read on input.
- Execution errors terminate in a post-mortem dump, providing a complete execution history that includes procedure invocations, variable values, type of error, etc
- the use of separately-compiled procedures in Pascal, FORTRAN or other languages is supported by the linkage-edit version. Thus one can build up a library of pascal procedures or use a pre-existing library of FORTRAN routines.

3. Measurements

- compilation speed about 2,500 chars/sec on an IBM 360/65
- compilation space : 128 K for small programs

160K for medium programs
220K for the compiler

- execution speed : comparable with Fortran $G$, at times better than FORTRAN H.
- execution space : about 30K plus the size of the compiled code, stack and heap
Compiled code is fairly compact - the compiler itself occupies 88k.

9. Reliability.

The system was first distributed in its current form early in 1978. It is currently used at about 90 sites. Reliability reports have been generally good to excellent.

## 10. Development Method

The compiler was developed from Nageli's trunk compiler and bootstrapped using Pascal-P by Hikita and Ishihata, who got it running on a HITAC-8000 computer (similar instruction set to IBM360). This version was further developed by Tobias and Cox for use under the OS family of operating systems on IBM360/370 computers. The compiler is written in Pascal 8000 ( 6000 lines) and runtime support is in 360 Assembler ( 3500 lines). Cox and Tobias spent about 10 person-months on the system ost of this time was spent improving the os support and adding enhancements to what was already a very workable system
11. Library Support.

The linkage-edit version has the ability to perform separate compilation of procedures or functions. These can be stored in a library and selected by the linkage editor as necessary. It can also link to routines written in FORTRAN or other languages which use a FORTRAN calling sequence. To use an externally compiled
routine, one must include a declaration for it. Such declarations consist of the procedure or function skeleton followed by the word 'pascal' or 'fortran'. The linkage-editor then automatically searches for that routine when it is linking the program. Global variables are accessible to externally compiled Pascal routines. pascal procedures cannot be overlayed.

A symbolic dump of local variables and traceback of procedures called is provided on detection of execution exrors.
12. Future Developments.

Version 2.0 is currently under development.
IBM 360/370 Berlin

We have received no new information on this VM370 ( $C P+$ CMS ) and os implementation since that which we published last year in Pascal News issue: \#11: 99-100.

IBM 360/370 Grenoble

We have received no new information on this OS/MVT and VS/MFT implementation since that which we published last year in Pascal News issue: $\# 9-10: 100$.

## IBM 370 London

We have received no new information on this CMS implementation since that which we published last year in Pascal News issue: \#11: $96-98$.

IBM 360/370 Manitoba

We have received no new information on this MFT, MVT, VS1, VS2, MVS, and CM implementation since that which we published last year in Pascal News issue: Mg

## STANFORD UNIVERSITY

Stanford Linear Accelerator Center<br>Sept. 15, 1978

Nail Address
SLAC, P. O. Box 4349 Stanford, California 9430

Dear Andy:
This is to announce the release of a new version of the Stanford PASCAL Compiler. This version provides comprehensive runtine checking as well as provisions for user-requested or post-mortem (symbolic) dump, separate compilation and generation of progran profile (i.e. frequency of execution of source and, except for a few restrictions, about $50 ø \emptyset$ lines long described in Jensen \& Wirth's "User Manual and Report". There are also some minor extensions to allow timing and clean termination of prograns without GOTOs across procedure boundaries.

The postprocessor, which translates the output of the compiler into IBM/37D assembly or object code, has also grown to source lines but the compilation/postprocessing time for
the compiler has renained almost unchanged (i.e. about il seconds of compilation followed by 5 seconds of postprocessing on the $370-168$, or a compilation rate of 590 lines per second). The combined system is still capable of self compiling in a 128 K region, but a larger area improves the I/O efficiency by allocating larger buffers.

Our earlier decision in leaving the compiler as machine independent as possible and writing a separate program to explained in the pascal Newletter \#8) proved to be very helpful in simplifying the task of bootstraping the compiler on a set of drasticaly different target machines. For example, after analyzing the static and dynamic properties of programs expressed in the intermediate form, we concluded that this form was quite suitable for a very compact encoding.

A postprocessor, intended primarily for microprocessor environments, translates the full compiler into a mere $2 ø \mathrm{~K}$ bytes which could be run interpretively, or implemented by a processors. Another interesting outcome of this implementation was that a very small ( 3 K bytes) $8080 / 280$ based interpreter, in conjunctin with the obove postprocessor, resulted in microprocessor-resident compiler with a compilation speed o about loø times slower than the $379-168$ in terms of the $C P$

Independent from these justifications, there are also some ther projects involved in writing machine independent p-code optimizers which would potentially benefit all the programs which are translated into the common intermediate form before being tied to the final target machine.

In conclusion, the PASCAL p-compiler seems to have helped spread the use of PASCAL far more than the sophisticated (and The zurich group should be credited for its farsightedness in developing this compiler as a separate program as well as devining the original 'p, pseudo machine which has since established the common grounds for the portability of PASCAL systems.

Sincerely
S. He.2.g.tat

Sassan Hazeghi
Computation Research Group
P.S. The new version of the $37 \varnothing$ Compiler is available through HARE Program Library as well as Argonne. Code Center, the incroproces
Code Center

We have received no new information on this OS implementation since that which we published last year in Pascal News issue: \#9-10: 98-99.

IBM 360, 370 (Vancouver)
0. DATE/VERSION. Barry Pollack reported (* 78/8/7 *): "Pascal/UBC is almost ready for its next round of distributions--it is an upward compatible superset of the old Pascal/UBC system, which is upwards compatible with Standard Pascal. The system runs on IBM 360/370 and Amdahl 470 machines. We plan to begin this round of distribution in sept. or Oct.--of
course, the old system is still available."

1. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER. Barry W. Pollack and Robert A. Fraley, Department of Computer Science, University of British Columbia, Vancouver, British Columbia, Canada V6I IW5 (604/228-6794 or 604/228-3061).
2. Machine. IbM 370/168.
3. SYSTEM CONFIGURATION. The current version runs under the MTS (Michigan Time Sharing) operating system. The monitor may be modified with minimal effort to run under VS, OS, store. Division of the compiler into overlays for non-VM systems would be possible.
4. DISTRIBUTION. The current version is available for distnibution now, via 9 track magnetic tape. Costs will be limited to postage (and tape purchase, if one is not supplied).
5. DOCUMENTATION. A User's Guide describes completely the implementation's departures from the Jensen and Wirth Pascal User Manual and Report. (* Apparently not machine retrievable. *)
6. MAINTENANCE. No policy has been decided. It is anticipated that periodic upgrades and modifications will be distributed at least once a year. Reported bugs will be corrected as quickly as possible with notification to users.
7. STANDARD. The compiler provides numerous extensions and a few restrictions. A compiler option issues error messages when non-standard features are used. A complete description is contained within the documentation provided. A summary of the differences follows

## xtensions:

Strings are padded on the right with blanks
a case default label: "<>".
Optional ";" allowed before else.
The character eol has been retained.
packed is ignored.
Input of character strings using read is allowed.
Support of EBCDIC characters "\&", "|", and (logical not sign). (* Sorry, we use
"ASCII at Pascal News. *)
value section exists for variable initialization.
Hexadecimal integers are supported.
return code is available in pre-declared variable rcode.
FORTRAN subroutines may be called. [©irect access files are supported.
解
reset and rewrite, and insert for data packing
Restrictions:
Sets are limited to 32 elements (0..31).

Program heading is not used
Files may not be components of other structures.
et constructors may not include <expression>..<expression>.
nput is initially eol instead of the first character of the file. This is transparent when read is used
$\frac{\text { Projected }}{\text { McGarthy if. }}$
rand and lower precedence than relations
"Usual" precedence used throughout
Sets over the range $0 . .255$.
Better control of input and output formats.
8. MEASUREMENTS. The compiler is written in Pascal and is modeled after the CDC 6000 implementation, but it has been extensively modified and improved. The translator consists of approximately 8000 lines of Pascal code. The run-time library consists of approximately 500 lines of Pascal code. The ming has not been determined, but it seems faster than our Algol-W compiler. The code produced has been timed against Algol-W code and is almost uniformly $10-15 \%$ better. This is especially true of any program using a large number of procedure calls. The compiler compiles itself in less than 60 seconds of $370 / 168$ processor time. The compiler requires 320K bytes of core
9. Reliability. To date has been excellent. A student version of the translator has been running since September, 1976, with only one detected compiler error. The main system version has been in operation since December, 1975. All problems which have been encountered to date have ben corrected. (* Number of sites using system not reported. *)
10. DEVELOPMENT METHOD. The original translator was developed by Wirth and several graduate students at Stanford University as a partial re-write of the CDC 6400 version in 1972. The current translator and monitor have been extensively modified, a run-time library has been implemented, and a post-mortem symbolic dump package has been developed. The translator has been under continuous development at UBC since December, 1975, by two faculty members and one (* anonymous? *) graduate student
11. LIBRARY SUPPORT. Fortran routines can be called. The compiler generates standard os object modules.

## IBM 360/370 Williamsburg

We have received no new information on this OS/VS implementation since that which we published last year in Pascal News issue: \#11: 95-96.

## IBM Series I (East Providence

It has been reported that SPAN Management Systems; Westminister Industrial Park; East Providence, RI 02914; 401/438-2200 has developed a dialect of Pascal which they call TSS and which will run on the IBM Series 1 computer; but we have received no information from them on their system.

## IBM Series 1 (Reston)

We have received no new information on this implementation since that which we published last year in Pascal News issue: \#9-10: 85.

ICL -- Introduction

PCHICL - the Pascal Clearing House for ICL machines - exists for the purposes of :

- Exchange of Library routines;
- Avoidance of duplication of effort in provision of new facilities;
- Circulation of user and other documentation;
- Circulation of bug reports and fixes;

Acting as a "User Group" to pascal users and implementors;
There are currently about 40 people on PCHICL's mailing list, mainly in Computer Science departments and Computing Centres of U.K. Universities and Polytechnics. Any user of Pascal on ICL machines whose institution is not already a member of PCHICL should contact:

Hull College of Higher Education
Inglenure Avenue
Hull HU6 7LJ
England, U.K
(0482) 42157
All ICL Pascal users are urged to notify David of any bugs they find, any compiler modifications they make, any useful programs or routines or documentation they have written, anything they have that may be of use or interest to other users.

Pascal Compilexs for the ICI 1900 series (\& ICL 2903/4) D。A.Joslin,
May 22nd 1978

## 1. \#PASQ Issue 3

This compiler is the most suitable for ICL 1900s operating under George 4 , and for those with a large core store ( 256 K , say) operating under George 3. This is the compiler described in the Implementation Checklist in "Pascal News". It incorporates a Diagnostics Package (written by D.Watts \& W.Findlay of Glasgow University) and a Source Library facility. It tekes 44K to compile most programs ( 60 K to compile itself). It may be obtained by sending a mag.tape (7-track NRZI 556 bpi or 9 -track PE 1600 bpi ) to the implementor, viz: Dr. J.Welsh,

Dept. of Computer Science, Queen's University, BELFAST, N.Ireland, BT7 INN.

## 2. \#PASQ Mark 2A

This compiler is suitable for all ICL 1900s (except 1901, 1901A, 1902, 1903: 1904, 1905) \& 2903/4s with at least 48 K of core: it is the most suitable compiler for ICL 1900s operating under George 2, and for those operating under George 3 where core is at a premium. The language processed (the language of the revised report) is identical to that procesied bJ H Haśa Issue 3, the compinex uescribed in the Impienentation Cheoklist in "Pascal News", but there is no Diagnostics Package or Source Library facility. The compiles takes 36 K to compile many programs, 40 K
to compile all but the most complex ( 48 K to compile itself). It was implemented orişinally by Queen's University, Belfast, and has been enhanced to include:

Selective compilation listing and insertion of mun-time checks;
Nesteà comments;
Improved compilation listing layout, and full text of compilation error messages;
Improved execution exror handling;
More efficient mathematical standard functions;
Facility to compile l5Ail procrams;
Specification of object-probran card \& line lengths;
Correction of various errors.
It may be obtained by sending a mag.tape (7-track NRZI 556 bpi or 9-track PE 1600 bpi ) to: D.A.Joslin,
(* address on previous page *)

## 3. \#XPAC HErk 1B

This compiler is suitable for all ICL 1900s 8: 2903/4s with at least 32 K of core. The language processed is Pascel Mark $I$, ie the language of the original report. The compiler takes 24 K to compile most prosrams ( 32 K to compile itself). It may be obtained by sending a mag.tape to Sussex (as in para 2 abcire).

## 4. Pascal-P

A Pascal to $P(4)$-code translator, configured for ICL 1900s : 2903/4s, may be obtained by sending a mag.tape to Sussex (as in para 2 above). This Is suitable for all ICL 1900s (except 1901, 1901A, 1902, 1903, 1904, 1905) \& 2903/4s with at least 32 K of core. The language processed is broadly the language of the revised report - see the Pascal-P section of "Pascal News". The translator takes 24 K to compile most programs (28K to compile itself). To complete the compilation process, either a p-code interpreter (based on the model interpreter provided) or a P-code to machine-code translator must be written.

## 5. Future Developrent

A two-stage Pascal compilex, which will be suitable for all ICL l900s (except perhaps 1901, 1901A, 1902, 1903, 1904, 1905) \& 2903/4s with at least 32 K or core, is to be produceã by Beirisst, possj.bly by October 1978. mhe language processed will be identical to that processed by $\# P A S Q$, and a Diagnostics Package and Source Library facility (George 3/4 only) may be provided.

ICL 1900 (Belfast)
0. DATE/VERSION. Checklist last updated 77/11/4.

1. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER. Jim Welsh, Colum Quinn, and Kath1een McShane, Department of Computer Science, Queens University, Belfast BT7 1NN, Northern Ireland, U.K. ( $*$ No phone number provided. *) Enhancements by David Watts and Bill Findlay, Computer
(* No phone Science Department, University of Glasgow, Glasgow G12 8QQ, Scotland, U.K. ( number provided. *)
2. MACHINE. ICL 1900 Series.
3. SYSTEM CONFIGURATION. Has been installed under George 3, George 4, Executive, MAXIMOP, and COOP operating systems. Requires 36 K ; uses CR, DA, LP files. (Source library facility only possible, and diagnostics package only practicable under George 3 or 4.)
4. DISTRIBUTION. Free--send 9 -track 1600 bpi PE or 7 -track 556 bpi NRZI tape to Belfast.
5. DOCUMENTATION. Belfast's Users' Guide (supplement to Pascal User $\frac{\text { Manual }}{\text { is distributed }}$ 解 $\frac{\text { Report }}{\text { the compiler. }}$ (Revised edition)) and implementation documentation is distributed with the compiler, Glasgow s Supplement of Computer Science, University of Glasgow, Glasgow, Scotland, G12 8QQ, United Kingdom (who also produced the Diagnostics package).
6. MAINTENANCE. No formal committment to maintenance. No plans for development in near future. Send bug evidence to Belfast, and also a note of the bug to PCHICL (see notice under ICL--Introduction) who circulate the bug reports and fixes to their members.
7. STANDARD. The level of the Revised Report; with:

Exceptions: There are no anonymous tag fields; files cannot be assigned, passed as value parameters, or occur as components of any structured type; Predefined procedures and functions cannot be passed as actual parameters; The correct execution of programs which include functions with side effect is not guaranteed; Only the first 8 characters of identifiers are significant; sets are limited to $x \cdot \cdot y$ where $0<=\operatorname{ord}(x)<=\operatorname{ord}(y)<=47$; The ICL 64 character graphic set is used for type char; packed is implemented, and text $=$ packed file of char; alfa $=$ packed array $[1 . .8]$ of char.
aditional predefined functions and procedures include: DATE, TIME, MILL, HALT, CARD; procedures ICL, ADDRESSOF allow use of inline machine code.
8. MEASUREMENTS. Compares favorably to Fortran, requiring about 32 K to compile. Generate code is better than that produced by the old 1900 Pascal compiler. (* Compilation speed not reported. *) Performance is better than most other ICL 1900 language processors (exceptions are incore compile-and-go batch systems of the WATFOR type).
9. RELIABILITY. Reported to be good. The compiler is in use at about 50 sites. (* Date first released not reported. *)
10. DEVELOPMENT METHOD. This compiler resulted from a complete rewrite of the old ICL 1900 compiler, which was bootstrapped from the CDC 6000 Zurich compiler. The new compiler is designed for portability, with a clean separation between semantic analysis and code generation. The compiler is about 14,000 lines of Pascal plus about 3500 lines of assembler code and produces absolute binary machine code. The post-mortem analysis program is about 2500 lines of Pascal.
11. LIBRARY SUPPORT. Allows access to Fortran routines.

ICL 2900 (Southampton)
0. DATE/VERSION. Checklist last updated 77/11/4.

Computer Studies Group; Faculty of Mathematical Studies; The University; Southampton, 509 5NH; England, J. K. OTOJ/J59122 x2270. Implementors: J. J. M. Reynolds; Computer Centr Queen Mary College; University of London; Mile End Rd.; London, El 4NS; England, U. K.; 01 9804811 x778 and H. J. Zell (deceased). The Pascal compiler will be distributed as a standard ICL program product. Contact the nearest ICL sales office or the Project Supervisor above.
2. MACHINE. ICL $2960,2970,2980$ series.
3. SYSTEM CONFIGURATION. VME/B and VME/K.
4. DISTRIbUTION. Gontact the nearest ICL sales office or the Project Supervisor above. (* No information provided on cost, tape formats, etc. *)
5. DOCUMENTATION. Standard ICL manuals will be available: a) Pascal Language Manual: operating system independant aspects of the Pascal. language. b) running Pascal Programs on $\mathrm{VME} / B$ and $\mathrm{VME} / \mathrm{K}$ : information on how to run Pascal under the operating system.
6. MAINTENANCE. Full maintenance will be provided by the implementation group and/or ICL while the compiler is offered as an ICL product. The usual ICL procedure for bug reports will be adopted.
7. STANDARD. The compiler implements "all" [sic] features of the language as described in Pascal: User Manual and Report.
8. MEASUREMENTS. Code generated is fairly compact, the compiler itself producing 80000 bytes. This is better than the 2900 standard compilers. The (CDC) Pascal 6000 compiler compiles the 2900 compiler on a CDC 6400 in 82 seconds. The ICL compiler self-compiles on John Reynolds tells us, "I've determined that almost all time required for a compilation on the 2900 is just burnt up by the system and that hardly any time at all goes in the actual act of code generation." (* 77/7/8*) (* Execution speed of generated code not reported. *) The source listing is approximately 10,000 lines of Pascal and produces 80 k bytes of code. Approximately 160 k bytes of store are required to compile the
compiler.
9. RELIABILITY. The compiler has been extensively tested and seems to work fairly well. Current (* 77/12 *) reliability is moderate to good. (* Date of first release and number of sites using system not reported. *)
10. DEVELOPMENT METHOD. The compiler is written in Pascal and produces Object Module Format (OMF) compatible with all standard ICL compilers. The OMF module may be directly loaded or linked with other OMF modules. The compiler was bootstrapped using the 1900 compiler from Queen's University of Belfast as a base. Twenty-four person-months of effort
11. LIBRARY SUPPORT. As the compiler produces OMF modules, separate compilation and the inclusion of external procedures will be possible providing the necessary operating system facilities are present.

## IMSAI VDP-40

See Intel 8080

Intel 8080 Ann Arbor
We have received no new information on this implementation since that which we published last year in Pascal News issue: \#12: 64-66.

Intel 8080 INSITE
We have received no new information on this implementation since that which we ublished last year in Pascal News issue: \#9-10: 102.

## Intel 8080 (Minneapolis)

A 25-page report on "Tiny Pascal", a cross-compiler for a greatly restricted variant of Standard Pascal which is written in CDC or the Intel 8080 is avil 227 en niversity Comp 5545 ( $612 / 373-4181$ ) (6neapolis, $\mathbb{N} 55455$ (612/373-4181)

## Intel 8080 Munich

We have received no new information on this implementation since that which we published last year in Pascal News issue: $\# 12: 66$.

## Intel 8080 Stanford

We have heard reports that there is an implementation of Pascal for the Intel 8080 microprocessor that has been developed at Stanford University (Stanford Linear Accelerato 360/370 Stanford in this issue.

Interdata 7/16 San Diego

We have received no new information on this implementation since that which we published last year in Pascal News issue: $\# 12: 67$.

Interdata $8 / 32$ Manhattan, Kansas
We have received no new information on this implementation since that which we published last year in Pascal News issue: $\# 9-10$ : 103-104.

Interdata $8 / 32$ San Diego

We have received a copy of two reports ( $*$ dated $78 / 5 / 2 *$ ) on cross-compilers for Sequential and Concurrent Pascal which run on the Univac 1100 series and produce code for the Inter 1.1 714/225-2366. Intel 8080a UCSD

See DEC LSI-11 UCSD

See IBM 360/370

Marinchip Systems M9900

## Marinchip Systems

omputer hardware and software

$$
16 \text { Saint Jude Road }
$$

Mill Valley, Ca. 9494
(415) 383-1545

## Marinchip 9900 Sequential Pasca:

Implementation Checklis

1. Distributor/Implementor/Maintainer

> John Walkex Marinchip Systems 16 St. Jude Road

Mill Valley, CA 94941 (415) 383-1545
2. Machine

Texas Instruments TMS9900. This system runs on the M9900 CPU
which adapts the TMS 9900 to the S-100 (Altair/IMSAI/etc.) bus.
3. System configuration.

Runs under Marinchip Disc Executive. Minimum configuration to compile compiler is 56 K bytes main memory and one IBM-compatible floppy disc drive
4. Distribution

Pascal is available to purchasers of the M9900 CPU board for $\$ 150$. the system is distributed on an IBM-compatible floppy disc in Disc Executive format.
5. Documentation.

Documentation supplied is a supplement to Per Brinch Hansen's book, The Architecture of concurrent Programs, and his Sequential Pascal
6. Maintenance policy.

Bug reports accepted from purchasers of the system. Fixes are available at reproduction cost. System is brand new: no maintenance track record
7. Standard.

Based upon Per Brinch Hansen's Sequential Pascal, so all comments in the Pascal Variants section about that compiler apply to this one too The lexical scanner has been modified to permit identifiers to be upper and lower case (case does not affect matching), to accept curly equential Pascal syntax still accepted as before.
8. Measurements.

The M9900 permits use of either 8 bit memories or 16 bit memories
With 8 bit memories, the memory cycle time is 3 us, and with 16 bit memories, the cycle is 1 us. Which kind of memory is used has a radical effect on performance. With 8 bit memories the with 16 bit memories, the speed is about 130 characters per second. To good benchmarks have been run to judge execution speed. Based on the performance of the original PDP-11 system and comparison of the PDP-11 and 9900 interpreters, we expect performance to range between $25 \%$ and $50 \%$ of native machine speed based upon instruction mix.

Reliability.
No extensive testing of the system has been done by users. However since the compiler has been compiled through itself without problem the system is felt to be quite stable.
10. Development method.

The system was bootstrapped from the PDP11/45 version of Sequential ascal. The interpretive object code was loaded onto the 9900 system and an interpreter was written for the interpretive code. Rather than implement the entire Solo operating system with which the compiler alls on the Marinchip Disc Executive. The execution environment
a Sequential Pascal program is completely simulated. The compiler root segment and seven passes were then compiled through the compiler. The code interpreter and operating system interface total 3000 lines
of 9900 assembly code. The compiler was transported and brought up
in less than one man-month. The implementor has previously written and
moved numerous compilers, but this was the first work on Pascal.
11. Library support.
eparately-compiled Sequential Pascal programs may call each other,
passing up to 9 arguments of type INTEGER, BOOTEAN, POTNTER, or
DENTIFIER ( 12 character array of CHAR). The program is loaded
coresident with its caller, executed, and a completion status is
returned to the caller (termination type and source line). Program
calls may be recursive, and nesting depth is limited only by
available memory and a configuration parameter. A utility program
may be called either from the user terminal, or from another program.
12. General comments

The Sequential pascal compiler was found to be excellently documented ery reliable in our tests, and extremely easy to move. The current
9900 system is source and object compatible with the PDPIl version.
Efficiency considerations may force divergence from the current object code compatibility.
MITS Altair 680B

See Motorola 6800 St. Paul

Mitsubishi MELCOM 7700

We have received no new information on this implementation since that which we published last year in pascal News issue: \#9-10: 104-105

Stephen P．Smith；P．O．Box 841；Parksley，VA 23421；804／665－5090 is working on a Pascal system for the MOS Technology 6502 chip（using the Ohio Scientific Industries Challanger I system）．The system will originally be the minimum subset of Pascal needed to write its own compiler．The original version will cross－compile on any machine which
supports a full standard Pascal compiler．The compiler will then convert itself to 6502 supports a full standard Pascal compiler．The compiler will then convert itself to 6502
machine code and further revisions will then be written in the Pascal subset resident on machine code and further revisions will then be written in the Pascal subset resident on
the 6502 ．As of $77 / 12$ ，the parsing procedures were completed and undergoing testing on a DEC－10．

## MOS Technology 6502 UCSD

See DEC LSI－11 UCSD．

Motorola 6800 St．Paul

We have received no new information on this implementation since that which we published last year in Pascal News issues：\＃9－10：105．\＃11： 102.

Motorola 6800 UCSD

See DEC LSI－11 UCSD

Motorola 6809

See Motorola 68000

Motorola 68000

See also Motorola 6800.
Computer Weekly reported on 78／9／7：＂Giving further credence to the view that Pascal could become the dominant high－level language of microcomputing，Motorola Semiconductor has revealed that this sof tware will be the prime language supported by its new processor， MACS，due to be unveiled early next year．
＂As an intermediate upgrade to MACS，Motorola will also be offering Pascal on its existing 6809 processor chip．The language is already available for the 6800 family from an independant source．
ly next year，and to show its lineage with the 6800 is expected to see the light of day known as the 68000．＂

Nanodata QM－1 California
We have received no new information on this implementation since that which we published last year in Pascal News issue：$⿰ ⿰ 三 丨 ⿰ 丨 三 一 ⿻ 上 丨 又-10: 105$.

We have received no new information on this implementation since that which we published last year in Pascal News issue：$\# 9-10$ ： 105.

Norsk Data NORD－10 CERN

We have received no new information on this implementation since that which we published last year in Pascal News issue：\＃9－10： 106.
Norsk Data NORD－10 0slo

We have received no new information on this implementation since that which we published last year in Pascal News issue：\＃9－10： 106.

## North Star Horizon

O．DATE／VERSION．Summer 1978.
1．IMPLEMENTOR／DISTRIBUTOR／MAINTAINER．North Star Computers； 2547 Ninth St．；Berkeley，CA 94710；415／549－0858

2．MACHINE．North Star Horizon Z－80 based system．
3．SYSTEM CONFIGURATION．Requires 48 K of RAM and the Micro Disk System
4．DISTRIBUTION．\＄49 including software on diskette and complete documentation（＊not known if this is machine retrievable＊）．

5．DOCUMENTATION．（＊No information reported．＊）
6．MAINTENANCE．（＊No information reported．＊）
7．STANDARD．The system is an implementation of UCSD Pascal，which varies from Standard Pascal．
8．MEASUREMENTS．（＊No information provided．＊）
9．RELIABILITY．（＊No information provided．＊）
10．DEVELOPMENT METHOD．（＊No information provided．＊）
11．LIBRARY SUPPORT．（＊No information provided．＊）

## Northwest Microcomputer Systems 85／P

[^1] documentation．

## Ohio Scientific Industries Challanger I

See MOS Technology 6502.

Prime P-300 and P-400 Hull
THE UNIVERSITY OF HULL
Telephonc: Hull 46311
ULL hue 7rx. england

Department of Computer Studies
30th August, 1978
Dear Andy
We're enclosing a fuller set of notes for our implementation of PASCAL on a PRIME 300. The work is now almost complete and we're very pleased with the result.

We have appended some extra sections to the notes. One of these deals with other implementations on PRIMEs and provides a brief summary of the
information we hold on them. Unfortunately we can't do a comparison of ald implementations since the Georgia Tech. version only runs on a PRIME 400.

Thanks again for your work with "PASCAL News".

## Yours sincerely

lan

## Barry Cornelius.

Ian Thomas.
Dave Robson

## THE UNIVERSITY OF HULL'S PASCAL COMPILER

## FOR PRIME 300 COMPUTERS

. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER: Barry Cornelius, Ian Thomas o ave Robson; Department of Computer Studies, University of Hull, Hull, HU6 7RX England; Hull (O482) 497951
2. MACHINE: Developed on a PRIME 300 but will also run in 64 R mode on a PRIME 400.
3. SYSTEM CONFIGURATION: The PRIME 300 currently has 64 K words running under PRIMOS-3 Revision Io.
4. DISTRIBUTION: Two versions of the compiler have been released to PRIME (U.K.) for evaluation and testing. It is hoped to have a distribution arrangement agreed for evaluation and testing.
5. DOCUMENTATION: A 30 page manual describing the PASCAL system is available in machine-readable form. It includes instructions on how to build a new compile
6. MAINTENANCE POLICY: This will depend partly on the agreement with PRIME evertheless we intend to correct reported errors for the next few years.
7. STANDARD: The PASCAL-P variant of PASCAL is implemented.. Some of its restrictions have been removed and some extensions have been added. The xtensions include external procedures (see 11 below) and an initialisation facility for variables in the outermost block.
B. MEASUREMENTS: When range-checking code is produced the compilation speed is approximately 550 characters/second. When code with no checks is required the speed is speed (without that ame as fortran's compilation speed (without trace or checking)

PASCAL input/output is considerably superior to FORTRAN's input/output. A text copying program takes about 4 times longer to execute in FORTRAN than PASCAL.

We do not have any comparisons for processor-bound programs since no-one can be persuaded to write a sufficiently large program in FORTRAN! However, the code is currently performed.
9. RELIABILITY: The compiler is very reliable and will reach a stable state by September 1978. It is hoped that the first release will then be available. As stated in 4 above, a preliminary release of the compiler is currently avail able on PRIME (U.K.)'s demonstration machines.

The Run-time Support and the input/output routines have been designed so that, when an execution time error occurs, an error number is output together with a "wordcount". The wordcount is the address relative to the start of the appears at the start of each line of the corpil The value of the wordcount can be traced to the line of the source program at which the error occurred. 10. DEVELOPMENT METHOD: The code generation sections of the PASCAL-P compiler
have been extensively rewritten to generate 64R mode PMA. It is a true compile rather than a compiler/interpreter system or a threaded code interpreter. The compiler is now some 6000 lines and compiles itself (without a compilation listing) in 300 C.P.U. seconds on the configuration described in 3 above. The first version of the compiler was developed from the PASCAL-P compiler on the University's ICL 1904S using the Belfast Mk. 2 compiler.
11. LIBRARY SUPPORT: Calls of external procedures are permitted. The parameterpassing protocol is a superset of that used by PRIME's standard system routines.
2. OTHER IMPLEMENTATIONS: There are a number of other implementations of PASCAL on PRIME machines. Some of these are described in more detail in a article we wrote for the Bulletin of the European PRIME Users Association, (see "PASCAL", E.P.U.A. Bulletin, Volume 4, Issue 1 (June 1978)).
(i) Per Brinch Hansen's Sequential PASCAL - very slow.
(ii) University of Brunswick's PASCAL compiler. Translates into modified poode which is subsequently optimised and translated into relocatable binary. The code produced contains calls to routines to perform Pcode instructions and it is thus a threaded code system. Compilation takes approximately 3 to 4 times as long as the University of Hull's imple faster than the assembly of the PMA that our implementation produces.
(iii) Georgia Tech's PASCAL compiler. The compiler was developed for a PRIME 400. From "PASCAL News" \#12 the current version appears to be a threaded code interpreter.
3. FUTURE PLANS: It is likely that we will implement translation into relocaable binary in the near future. The additional compilation time overheads will probably be offset by the reduction in the amount of character input/output currently necessary to output PMA text.

More of the restrictions af the PASCAL-P subset are also likely to be removed. It is possible that we will implement the post-mortem dump facility removed. (in PASCAL) by Glasgow University for the ICL Belfast Mk. 2 compiler.

We have received no new information on this implementation since that which we published last year in Pascal News issues: \#9-10: 106. \#12: 67.

Processor Technology SOL

According to Ralph I. Palsson, Customer Applications Manager: Processor Technology Corp.; 7100 Johnson Industrial Dr.; Pleasonton, CA 94566; 415/829-2600: "We do not currently ( $* 78 / 1 / 11$ *) have any intentions of providing Pascal. We will be providing a FORTRAN compiler this spring as well as PILOT... Providing good software support for users
of Processor Technology hardware is one of our primary committments [sic]. As of this time, there has been relatively little demand for Pascal. Consequently our softwar emphasis has been in other areas.

According to S.M. Sokolow, Editor; Solus News; 1690 Woodside Rd. 219; Redwood City, CA 94061 (* 78/10/13 *): "We're in the process of preparing to distribute the Stanford

Radio Shack TRS-80
See also Zilog Z-80.
According to Hugh Matthias, Radio Shack, 205 NW 7th St., Fort Worth, TX 76101; Radi Shack does not intend to produce a Pascal system for the TRS 80 now or at any time in the fure. "Tt appears to be to [sic] costly-nembl" (* 77/11/19 *)

RCA Spectra 70
See Siemans 4004, 7000 and Univac 90/70.

## SEL 8600

Jim Gilbert; Systems Structuring Technology; 30436 N. Hampton Rd.; Laguna Niguel, CA 2677; 714/640-5222 (work); 714/495-6039 (home) reports (* 78/9/30 *): "I am the implementor of the SEL 8600 \& SEL 32 P2 Pascal mentioned in Pascal News \#4 [Co-implementor Michael] Richmond is with D.G.C. in Carolina last I knew. I am available on a contract basis for language consulting."

SEMS T1600 Nancy, France

We have received no new information on this implementation since that which we published last year in Pascal News issue: $\# 9-10: 106$.

## Siemens 150 and 330

We have received no new information on this implementation since that which we published last year in pascal News issue: $\# 9-10$ : 107-108.
O. DATE/VERSION: 78/10/01 Version 2.0

For version 1 see \# $2 / 10$ : 108

1. Distributor/Implementor/Maintainer

Dr. M. Sommer
SIEMENS AG Dep: D AP GE 1
Otto-Hahn-Ring 6
D - 8000 München 83
Germany
2. Machine:

SIEMENS series 4004 and series 7000
3. System configuration:
all systems under operating system BS2000 (>= rel.3.0)
4. Distribution

- please contact implementor

5. Documentation

Machine retrievable user manual
6. Maintenance Policy

- please contact implementor -

7. Standard.

Standard PASCAL is accepted
Extension: Sets of any range (maxelements: 2048)
are implemented by minimal byte-strings, separate
compilation of PASCAL, FORTRAN,-procedures and PASCAL-

Modules many additional standard procs.

- compiler cptions, like optimise, xref, debug, codelist, etc.
- compiler instructions like copy from include-lib, skip.

8. Measurements: (For a SIEMENS 7.755)

Compilation speed: 3200 chars (incl. blanks) /second 140 lines/second
(*speed is depending on options/listings*)
Execution speed and execution space of an
average of 6 test programs including prim, queens,
palindromes, quicksort. etc.

| PASCAL version 2 |  | PASCAL version 1 | other language |
| :--- | :---: | :---: | :---: |
| SPACE (bytes | 326 | 580 | 446 |
| TIME (sec) | 4.2 | 7.8 | 5.2 |

9. Reliabiiity
is hoped to be excellent as the reliability of version 1
is excellent. (Used by ca. 30 sites.)
10. Development method

Developed from version 1 (developed from PASCAL P).
New codegenerator - and other extensions.
Length is ca. 14000 lines of compact PASCAL.
Effort (Version $1 \rightarrow$ Version 2) ca. 20 MM.
11. Library support
a) Standard modules from Assembler, Fortran, Cobol
b) Procedures and modules written in PASCAL

Standard Linkage. Copy from libraries in Source.
SOLAR 16-05/40/65
See SEMS T1600.

SouthWest Technical Products

See Motorola 6800.

Telefunken TR-440

TERAK 8510, 8510A UCSD

See DEC LSI-11 UCSD.

Texas Instruments TI-ASC
The only new information we have received on this implementation since that which we published last year in Pascal News issue: $\# 9-10: 109$ is the rumor that the implementation is being done by the Advanced Software Technology group and is currently ( $* 78 / 2 / 28$ *) i the debugging stage; but that the compiler is probably not intended for outside istribution.

Texas Instruments TI-980a
George Cohn, Wrubel Computer Center, Indiana University/HPER, Bloomington, IN 47401 (812) 337-1911, has had a Pascal version running for quite some time in the Computer Science Department, although no formal distribution arrangements have been made.


0. DATE/VERSION

Release 1.4.0, May 1978.

1. DISTRIBUTER/IMPLEMENTATION/MAINTAINER

Implemented by Texas Instruments. Information is available from TI sales offices, or write to:
Texas Instruments
Digital Systems Division, MS 784
P. O. Box 1444
Houston, Texas 77001 Houston, Texas 77001
or call (512) 258-7305. Problems should be reported to:
Texas Instruments
Software Sustaining, MS 2188
P. O. Box 2909

Texas 78769
or call (512) 258-7407.
2. MAChine

TI 990/10
3. SYSTEM CONFIGURATION Runs under the DX10 operating system (release 3) on a TI
DS990 Model 4 or larger system, which includes a $990 / 10$ with 128K bytes of memory and a 10 megabyte disk.
4. Distribution

Available on 9-track magnetic tape (either 800 or 1600 bpi ) or on a disk pack for a TI miodel DS10, DS31, DS25, or DS50 disk drive. Contact a TI salesman for a price quotation.
5. DOCUMENTATION

Complete user-level documentation is given in the "TI Pascal User's Manual", TI part number 946290-9701.
6. MAINTENANCE POLICY

TI Pascal is a fully supported product. Bug reports are welcomed and maintainence and further development work are in progress.
7. . STANDARD

TI Pascal conforms to "standard" Pascal, with the following principal exceptions:

* Functions cannot alter global variables
* A GOTO cannot be used to jump out of a procedure.
* The control variable of a FOR statement is local
* The pecedance of Boolean operators has been
modified to be the same as in Algol and Fortran.
* The standard procedures GET and PUT have been
replaced by generalized READ and WRITE procedures.
TI Pascal has a number of extensions to standard Pascal, including random access files, dynamic arrays, ESCAPE and ASSERT statements, optional OTHERWISE clause on CASE statements, and formatted READ.

8. MEASUREMENTS

The compiler occupies a 64 K byte memory region. Compilation speeds are comparable to the 990 Fortran compiler.
9. RELIABILITY

There are some known problems which are currently being
worked on, but none are so serious that they can't be worked
around. The system has been used by several different groups within TI since October of 1977, and by a number of outside customers since May of 1978.
10. DEVELOPMENT METHOD

The compiler produces object code which is link-edited with run-time support routines to form a directly executable program The compiler is written in Pascal and is self-compiling.
11. LIBRARY SUPPORT

TI Pascal supports separate compilation of routines and
allows linking with routines written in Fortran or assembly
language.

Texas Instruments 9900/4 Vienna

We have received no new information on this implementation since that which we published last year in Pascal News issue: 非-10: 109.

Univac 90/30

We have received no new information on this implementation since that which we published last year in Pascal News issue: \#9-10: 109.

Univac 90/70

See Siemens 4004, 7000 series.

## UNIVERSITY of PENNSTLVANIA philadelphia 19104

## The Moore School of Electrical Engineering D2

 Department of Computer and Information ScienceApril 20, 1978
Dear Andy,
I just wanted to let you know about the PASCAL 8000 implementation which I recently brought up on our Univac $90 / 70$ (VS/9 operating system).

The system is based on the Australian AEC compiler of Cox, Tobias, Hikita and Ishihata (which is quite an excellent piece of software), and was implemented by modifying the runtime system to interface with VS/9. Only the compile-and-go version has been implemented at this time. All features of the Australian compiler have been retained, and additional support added for some VS/9 features: the system files SYSDTA, SYSLST, SYSOUT, SYSIPT, SYSOPT and * are supported, and a COMMAND function has COMMAND ('/ERASE filename'); This implementation will probably not run on Univac Series 70 VMOS without modification, since interrupt handling is done with operating system features that I am told are specific to vs/9.

No formal distribution plans have been made, but anyone who is interested (hopefully with software to trade) should contact me at P.O. Box 8191, Philade1phia PA


## Univac 1100 (Copenhagen)

We have received a copy of a. 60 page users manual (* dated 77/8*) titled "A Pascal Compiler for the Univac 1100 Series" which is available from the implementor.
0. DATE/VERSION. Checklist not updated since 77/08.

1. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER. J. Steensgaard-Madsen, DIKU (Datalogisk Institut Kobenhavns Universitet), Sigurdsgade 41, DK-2200 Copenhagen N., Denmark. (* No phone number reported. *)
2. MACHINE. Univac 1100 series
3. SYSTEM CONFIGURATION. Exec-8 operating system. (* Minimum hardware requirements not reported. *)
4. DISTRIBUTION. The charge for distribution from Datalogisk Institut is Dkr. 200. The distributors are attempting to maintain a distribution tree to reduce costs and hassles. purchasers must sign a license agreement. The system is released only in relocatable form.
5. DOCUMENTATION. A 19-page machine-retrievable supplement to the $\frac{\text { Pascal }}{}$ User $\frac{\text { Manual }}{\text { and }}$
Report is available. It is "A Pascal Compiler for the Univac 1100 machines", by $\frac{\mathrm{J}}{\mathrm{J}}$. Report is available. It is A Pascal Comp
Steensgaard-Madsen and Henrik Snog of DIKU.
6. MAINTENANCE. There is no promise of maintenance, but bug reports are required unde the license aggreement.
7. Standard.

Deviations from the standard: Reset(f) on any textfile $f$ will cause eof $(f)=$ fals and $e o l n(f)=$ true; Parameter types of formal procedures and functions must be specified. and eoln $(\mathrm{f})=$ true; Parameter types of formal procedures and functions must be specified. actual parameters.

Extensions: otherwise in case statements; conformant array parameters.
Machine dependencies: Sets may have 72 elements, char is defined as ( 6 -bit) Fieldata ASCII is an additional type; real is double precision always
8. MEASUREMENTS. Compilation space is roughly 42 K ; speed is 100 lines per SUP second Comp iled programs run efficiently compared to other processors.
9. RELIABILITY. Excellent. (* Date first released and number of sites using system not reported. *)
10. Development method. Pascal-P with a team of 4 persons. '(* Person-hours to develop system not reported. *)
11. LIBRARY SUPPORT. External procedures may be written in Pascal or (ASCII) Fortran. Inclusion of assembler code is possible.

Univac 1100 Madison, Wisconsin

We have received no new information on this implementation since that which we published last year in Pascal News issues: \#9-10: 110-112. \#11: 103.

Univac 1100 (San Diego)

We have received a 33-page report on this implementation titled "Pascal 1100" which is available from the implementor.
0. DATE/VERSION. Checklist not updated since $77 / 08$

1. IMPLEMENTOR/DISTRIBUTOR/MAINTAINER. Michael S. Ball; code 632; Naval Ocean Systems Center; San Diego, CA 92152; 714/225-2366.
2. MACHINE. Univac 1100 Series.
3. SYSTEM CONFIGURATION. Exec-8 operating system; can be run in Demand mode.
4. DISTRIBUTION. As a member of USF, you may request a copy from Mike by sending a mag tape and noting any restrictions on it $s$ format.
5. DOCUMENTATION. A machine-retrievable supplement to the Pascal User Manual and Report entitled "Pascal 1100 " documents the implementation.
6. MAINTENANCE. (* No information provided. *)
7. STANDARD. Restrictions: entry, processor, and univ are reserved words; standar procedures and functions may not be passed as actual parameters; file of file is no allowed. Sets may have at most 144 elements. The compiler accepts the full Ascil character set. A compiler option allows processing of Brinch Hansen Sequential Pascal programs.
8. MEASUREMENTS. The compiler compiles into 34 K words and requires 6 K words of library routines. (* Compilation speed not reported. *) Self-compilation requires about 15.5 K for stack and heap.

Execution times for code compiled by Pascal was compared with code generated by the NUALG and ASCII FORTRAN processors. Fortran's local optimization was taken as a bas Pascal compiler (Software - Practice and Wxpertence Vol Wirth s paper on he design . results are summarized in the following table.

9. RELIABILITY. Quite good; it should approach excellent. The system has been in local use since about February, 1976, and it has been installed at 25 sites (11 university, 4 government, 10 industry)
10. DEVELOPMENT METHOD. The compiler was developed from Pascal-P2. (* Person-hours to develop system not reported. *)
11. LIBRARY SUPPORT. Generated code can be linked to subprograms written in Fortran or assembler.

Varian $\mathrm{V}-70$ volce

We have received no new information on this implementation since that which we published last year in Pascal News issue: \#9-10: 112 .

Western Digital Newport Beach
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$\subset$ o a P o $A$
will begin in the first quarter of 1979.
Additionally, Western Digital offers a wide range of chip-level products which have been successfully used in a variety of applications including the following.

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- Terminals and Printers
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Flease call our regional offices or this author here at Newport Beach for additiona! information.

- Western - Mr. Ed Raether, Los Gatos, California (408) 354-28! 3
- Central - Mr. Dave Renwick, Trov, Michigan (313) 643-4482
- Eastern - Mr. Bob Green, Marblehead, Massachiusetts (617) 631-6466

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Sincerely,
$\vec{\alpha}$
J.T. Boren

Marketing Manager
Computer Froducis Division

Xerox Sigma 6, 9 Quebec

We have received no new information on this implementation since that which we co published last year in Pascal News issue: \#9-10: 112.

Thank you for your interest in Western Digital's innovative new Pascal microeng net product line. We are pleased to enclose our initial literature whith will soon be followed by more conclusive and detailed data sheets.

Our first product offerings for the Pascal MICROENGINE are at both the system and chip level. The desktop system (CP90007B-OX) configured in a stylized enclosure retails for \$2995, although a special introductory offer of $\$ 1995$ is in efrect for the tirst 500 customers to raserve a system. A 20 down payment must ascompany orders for this special cffer. Crders should be accompanied by the model number
(above) with the appropriate "-ox" suffix to specify the ciskette type for receipt of softivare: -03 and -04 for 9 irch standard diskette, sirigle and double densitv, respectively; ard -05 and -06 for $5 \mathrm{l} / 4$ inch mini disketro, single and double derisity, respectiveiy. The chif set (CP 90008B-02; retails for $\$ 195$. All pricas are sity, respectively. The chif ser
subject to applicable tax. Sct:' products are offered to the OEM and retail market segments with corresponding price schedules targeted to those markets. Deliveries

## Xerox Sigma 7 Toky

See also CII 10070 and CII IRIS 80. We have received no new information on this implementation since that which we published last year in Pascal News issue: \#9-10: 112.

Zilog Development System

See Zilog Z-80 UCSD.

Andy Mickel
University Computer Cente
227 Experimental Engineering Building
Minneapolis, Minnesota 55455

Dear Andy,

O\%er the past sevaral months I have been working on a PASCAL compiler for the Ziloz 7. Sij, The project is now at the point where it generates pretty good code. Any further enhancements will have to be done either at a much siower pace or by somebody else.

I started with the P4 compiler and wrore a PASCAL program that translated the PCODE to $Z 80$ assembly mnemonics. I discovered that for the sake of efficiency of space, many of the PCODE instructions have to generate calls to a next step was to eliminate the $z 80$ assembler. I next step was to eliminate the an intermediate file into my postprocessor so that it could generate standard 280 object code. My final step was to move the logic from the postprocessor up into the main compiler, making it one pass. In order to avoid having an intermediate file, it assembles the code for each procedure body in a chain of arrays allocated from the heap. Thus, the largest procedure determines the amount of memory needed by the compiler. The compiler itself still thinks it is zenerating PCODE: I have and a couple of others so that instead of printing PCODE mnemonics onto the object file, they call z80 code generating routines which do the real work. As a bonus, I now have actual $Z 80$ addresses alongside my listings, making breakpointing convenient.

My compiler is probably doomed to remain a cross compiler for the duration of its existence. Its output code cannot hope to compete with interpreter code for efficient takes up about 4 K bytes, and a good interpreter that used that package extensively could probably be written in a few hundred bytes. The compiler could be remodified to generate a binary interpreter code in $Z 80$ loader format. Then some really big programs could be run on the $Z S O$, perhaps even the compiler itself. But that is another project

I have made some little white modifications to the compiler to make it more convenient to the microprocessor programmer. These include the following: 1) Files. Not really an extension, since they are
part of the language definition. In the current
implementation, file names actually serve as device designators
2) Nondescriminated variants. Also standard PASCAL.
3) ASCII coding. CASE statement constants and SET constants based on characters are translated t ASCII. This is absolutely necessary for cross
4) External procedures and functions. These must be accompanied by an absolute address, since we have no relocating loader.
5) Hex and octal in the source code. 256, for example, is represented as 100 H for hex and 400 Q for octal.
6) Hex output on textfiles

Example: WRITE(OUTPUT,A: 4 HEX),
Under the present configuration, characters range from 0 to 255, sets may contain up to 128 elements, integers ar' I was once tempted to cut integers down to 16 bits, which would work wonders for efficiency, but Al Towell talked me out of it. What can you do with a 16 bit integer?

PASCAL can provide a convenient medium for applications systems with the help of assembly language procedures for tight spots. We seem to be getting into Z80s more and more code i worked on will see plenty of action on the fron lines.

I am sending a listing of my runtime support packag and listings of a couple of compilations for you to look over. Tell me what you think.

Respectfully,

George Cohn III

Zilog Z-80 UCSD
See also DEC LSI-11 UCSD.
We have received no new information on this implementation since that which we ublished last year in Pascal News issue: 非-10: 112 .

Zilog Z-8000
It is believed that many of the present Zilog Z-80 Pascal systems could be easily modified to run on the $\mathrm{z}-8000$; since Zilog says "using an automatic translator, present users of the $Z-80$ can easily convert to the $Z-8000$, since the $Z-8000$ instruction set is in effect, a superset of the $\mathrm{Z}-80$ instruction set." We would appreciate hearing from
anyone who has made such a conversion; or from anyone who has developed a Pascal syste directly for the $\mathrm{z}-8000$.



IBM 1130
IBM 303x
\＃9\＆10： 101
IBM 360， 370

\＃9\＆10：95－101．
\＃11：93－100．
\＃12： 64.
IBM Series 1
ICL $1900^{\circ}$
\＃9\＆10：101－102．

$$
\begin{aligned}
& \text { \#11: } 100-1 \\
& \# 13: 110 .
\end{aligned}
$$

ICL 2900
\＃9\＆10： 102.
 1 13： 111
IMSAI VDP－40
Intel 8080 Intel 8080.
Intel 8080，8080a \＃\＃\＆10：102－103．
$\begin{array}{ll}\text { \＃11：} & 102 . \\ \text { \＃12：} & 64-66 .\end{array}$
\＃12： $64-66$.
$\# 13: 112$.
Intel 8085
Intel 8086 Intel 8080.
See Intel 8080.
Interdata $7 / 16$
$\# 9610: 103$.
$\# 12: 67$.
Interdata $7 / 32,8 / 32$
\＃9 10 10： $103-104$.
$\# 12: 67$.
\＃12： 112.
ITEL AS／4，AS／5
See IBM 360， 370.
Marinchip Systems M9900
Mit tsub ishi MELCOM 7700
非\＆10：104－105
MITS Altair 680b
See Motorola 6800
MITS Altair 8800
MOS Technology 6502 See also DEC LSI－11． \＃13： 114.
Motorola 6800
\＃\＃9\＆10： 105
$\# 11: 102$.
Motorola 6809
see Motorola 68000.

| Motorola 68000 |
| :--- |
| 非13： |
| 114 |

Nanodata QM－1
\＃9910： 105.
NCR Century 200
NGR Century 200
Norsk Data NORD－10
非9 10： 106.
North Star Horizon
非13： 114.

Northwest Micro Systems 85／P
\＃12： 67.


\＃9810： 106.
Processor Technology SOL
Radio Shack TRS－8．
RCA Spectu： 116
See Siemens 4004， 7000
See Univac 90／70
SEMS T1600，SOLAR 16／05／40／65
\＃9\＆10： 106.
Siemens 150
See Siemens 330.
Siemens
330
\＃9\＆10：107－108．
Siemens 4004， 7000 \＃9\＆10： 108.

See Sems T1600．
Telefunken TR－440
$\# 9 \& 10: 108$.
\＃13： 117.
Terak 8510
Texas Instruments LI－11． nstruments
$\# \neq 9 \propto 10: 109$.
$\# 13: 117$
Texas Instruments TI－980a
Texas Instruments IIT－990， 9910
Texas Instruments $9900 / 4$

Univac $90 / 30$ ： 109
Univac $90 / 79810: 109$.
\＃9\＆10： 109.
\＃13： 118.
1100
Univac $\begin{aligned} & 1100 \\ & \text { \＃} ⿰ ⿰ 三 丨 ⿰ 丨 三\end{aligned}$ \＆ 10 ：109－112．

Varian $V-70$
\＃9\＆10： 112.
Western \＃13： 120.



\＃9\＆10： 112.
\＃11： 103.
\＃13： 120.
Zilog Z－8000 120

Purposes: Pascal User's Group (PUG) tries to promote the use of the programming Tanguage Pascal as well as the ideas behind Pascal. PUG members help out by sending information to Pascal News, the most important of which is about implementations (out of the necessity to spread the use of Pascal).

The increasing availability of Pascal makes it a viable alternative for software production and justifies its further use. We all strive to make using Pascal a respectable activity.

Membership: Anyone can join PUG: particularly the Pascal user, teacher, maintainer, implementor, distributor, or just plain fan. Memberships from libraries are also encouraged.

See the ALL PURPOSE COUPON for details.

FACTS ABOUT Pascal, THE PROGRAMMING LANGUAGE:
Pascal is a small, practical, and general purpose (but not all-purpose) programming language possessing algorithmic and data structures to aid systematic programming. Pascal was intended to be easy to learn and read by humans, and efficient to translate by computers.
Pascal has met these design goals and is being used quite widely and successfully for:

* teaching programming concepts
* developing reliable "production" software
* implementing software efficiently on today's machines
* writing portable software

Pascal is a leading language in computer science today and is being used increasingly in the world's computing industry to save energy and resources and increase productivity.

Pascal implementations exist for more than 62 different computer systems, and the number increases every month. The Implementation Notes section of Pascal News describes how to obtain them.

The standard reference and tutorial manual for Pascal is:
Pascal - User Manual and Report (Second, study edition) by Kathleen Jensen and Niklaus Wirth
Springer-Verlag Publishers: New York, Heidelberg, Berlin 1978 (corrected printing), 167 pages, paperback, \$6.90.

Introductory textbooks about Pascal are described in the Here and There Books section of Pascal News.

The programming language Pascal was named after the mathematician and religious fanatic Blaise Pascal (1623-1662). Pascal is not an acronym.

Pascal User's Group is each individual member's group. We currently have more than 2712 active members in more than 41 countries. This year Pascal News is averaging more than 120 pages per issue.


[^0]:    if currnargin < slofaill

[^1]:    Northwest Microcomputer Systems； 121 East Eleventh；Eugene，OR 97401；503／485－062 offers the Northwest $85 / \mathrm{P}$ ；a self－contained Intel 8085 based microcomputer which includes 2 double density full size Shugart floppy disks（1 Mbyte online），54K of 450 ns Static RAM （I／O，etc．in PROM），Hall effect typewriter keyboard with numeric pad and 29 user definable function keys， 24 line 80 character $12^{\prime \prime}(30 \mathrm{~cm})$ Video RAM display， 2 serial ports and 16 parallel ports．The basic system includes with the hardware the CP／M operating system and the Pascal system for $\$ 7,495$ ．The Pascal compiler／interpreter runs at 725 files，screen－oriented editior，interactive source linked debugger，full

