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**Systems Reference Library**

## **IBM 1130 Assembler Language**

This publication contains the information necessary to write programs in the IBM 1130 Assembler language. Included are rules for statement writing, mnemonic codes and descriptions of operands, and descriptions of the instructions used to control the Assembler program.

## PREFACE

This manual describes the IBM 1130 Assembler language and defines the programming rules. It is intended as reference material for the writing of an assembler source program and the accomplishment of the steps required to produce the resulting object program. For those without programming experience or a knowledge of the principles involved, the IBM publication, Introduction to IBM Data Processing Systems (Form F22-6517), is suggested as preliminary reading.

For those without experience involving different number systems, i. e., binary and hexadecimal, the publication IBM Student Text: Number Systems (Form C20-1618) is recommended.

The reader should also be familiar with the following: IBM 1130 Functional Characteristics (Form A26-5881) and IBM 1130 Computing System, Input/Output Units (Form A26-5890).

The assembler language is valid for the 1130 Disk Monitor Programming system and the 1130 Card/

Paper Tape Programming System. The operating procedures for the Monitor Assembler are described in the publication IBM 1130 Disk Monitor System, Version 2, Programming and Operator's Guide (Form C26-3717).

The operating procedures for the 1130 Card/Paper Tape Assembler are described in the publication, IBM 1130 Card/Paper Tape Programming System Operator's Guide (Form C26-3629).

## MACHINE REQUIREMENTS

The minimum machine configuration for assembling programs is as follows:

IBM 1131 Central Processing Unit, Model 1,  
with 4096 words of core storage  
IBM 1442 Card Read Punch, or IBM 1134 Paper  
Tape Reader and IBM 1055 Paper Tape Punch.

## Third Edition

This edition is a major revision of the previous edition (C26-5927-2) which is now obsolete. Information has been added that enables the user to program the additional I/O units available with Version 2 of the 1130 Disk Monitor System.

Significant changes or additions to the specifications contained in this publication will be reported in subsequent revisions or Technical Newsletters.

Requests for copies of IBM publications should be made to your IBM representative or to the IBM branch office serving your locality.

A form is provided at the back of this publication for reader's comments. If the form has been removed, comments may be addressed to IBM Corporation, Programming Publications, Department 232, San Jose, California 95114.

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## INTRODUCTION

The IBM 1130 Assembler language replaces binary instruction codes with mnemonic symbols and uses labels for other fields of an instruction. Other features, such as pseudo-operations, expand the programming facilities of machine language. Thus, the programmer has available, through an assembler language, all the flexibility and versatility of machine language, plus facilities that greatly reduce machine language programming effort.

### Symbolic Language

Symbolic language is the notation used by the programmer to write (code) the program. A program written in symbolic language is called a source program. It consists of systematically arranged mnemonic operation codes, special characters, addresses, and data, which symbolically describe the problem to be solved by the computer.

The use of symbolic language:

- Makes a program independent of absolute core locations, thus allowing programs and routines to be relocated and combined as desired.
- Allows subroutines that can be written independently and that cause no loss of efficiency in the final program.
- Permits instructions to be added to or deleted from a source program without the user having to reassign storage addresses.

### Assembler Program

The assembler program converts (assembles) a source program into a machine-language program. The conversion usually is one for one — that is, the assembler produces one machine-language instruction for each symbolic-language instruction.

The 1130 Disk Monitor Assembler is a two-pass assembler. The source program is read into core from the principal input device and written on the disk for use in pass 2. During the first pass the symbol table is generated. During the second pass the object

program is created in the system Working Storage and the listing, if requested, is produced.

The IBM 1130 Card/Paper Tape Assembler is a two-pass program. It is loaded into the computer and is followed by the first pass of the source program. During the first pass, the source statements are read and a symbol table is generated. During the second pass, the source program is read again and the object program and/or error indications are punched into the first 20 columns of each source card. If paper tape is used, the second pass results in the punching of a new tape that contains both source statements and corresponding object information. Both card and tape object programs must be compressed (via a Compressor Program supplied with the assembler) into a relocatable binary deck (or tape) before they can be loaded into core storage for execution. The output from the second pass is called the list deck (or tape) and can be used to obtain a program listing of source statements and corresponding object statements.

### Subroutines

A library of input/output, arithmetic, and functional subroutines is available for use with the IBM 1130 Assembler.

The user can incorporate any subroutine into his program by simply writing a call statement (CALL or LIBF, whichever is required), referring to the subroutine name. The assembler generates the linkage necessary to provide a path to the subroutine and a return path to the user's program. The ability to use subroutines simplifies programming and reduces the time required to write a program.

A description of available subroutines is contained in the publication IBM 1130 Subroutine Library (Form C26-5929).

## FEATURES OF THE ASSEMBLER

The significant features of the IBM 1130 Assembler are summarized below. More detailed explanations are given later in this manual.

Mnemonic Operation Codes. Mnemonic operation codes are used for all machine instructions instead

of the more cumbersome internal binary operation codes of the machine. For example, the Subtract instruction can be represented by the mnemonic, S, instead of the machine operation code, 10010.

Symbolic References to Storage Addresses. Instructions, data areas, and other program elements can be referred to by symbolic names or actual machine addresses and designations.

Renaming Symbols. A symbolic name can be equated to another symbol, so that both refer to the same storage location. This makes it possible for the same program item to be referred to by different names in different parts of the program.

Automatic Storage Assignment. The assembler assigns consecutive addresses to program elements as it encounters them. After processing each element, the assembler increments a counter by the number of words assigned to that element. This counter indicates the storage location available to the next element.

Relocatable Programs. The assembler can produce object programs in a relocatable format; that is, a format that enables programs to be loaded and executed at storage locations different from those assigned when the programs were assembled.

Convenient Data Representation. Constants can be specified as decimal digits, alphabetic characters, hexadecimal digits, and storage addresses. Conversion of the data into the appropriate machine format of the 1130 System is performed by the Assembler. Data can be in a form suitable for use in decimal integer, fixed-point or real arithmetic operations.

Program Listings. For every assembly, the user can obtain a program listing. This listing can be produced either off-line (Card/Paper Tape Assembler) or on-line during the assembly process (Disk Monitor Assembler).

Error Checking. Source programs are examined by the Assembler for errors arising from incorrect use of the language. Where an error is detected, a coded warning message appears in the program listing.

**MNEMONIC CONCEPT**

Symbolic programming may be defined as a method whereby names and symbols are used to write a program. The symbolic language includes a standard set of mnemonic operation codes. Mnemonic operation codes are easier to remember than machine language codes because they are usually abbreviations for actual instruction descriptions. For example:

<u>Description</u>	<u>Mnemonic</u>
Add	A
Execute I/O	XIO

Each IBM 1130 machine instruction has a corresponding mnemonic operation code. In addition, there are some mnemonic codes that assign storage and others that allow the user to exercise control over the assembly process.

**FORMAT OF STATEMENTS**

A source program consists of a sequence of statements. These statements can be written on a standard coding form (X26-5994) provided by IBM. The information on each line of the form (Figure 1) is punched into one card or paper tape record or entered from the keyboard. The first position on the form (21) corresponds to card column 21 or to the first character of the paper tape/keyboard record. Space is provided at the top of the coding form to identify the program; however, none of this information is punched into the statement cards. The first 20 columns of an assembler source card must be blank.

**NOTE:** Keyboard input is acceptable only with the Monitor 2 Programming System.

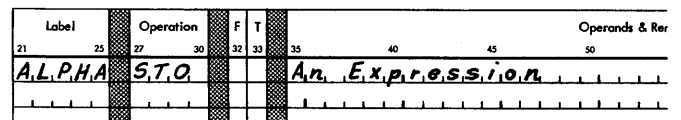
Statement Fields

An assembler statement is composed of one to seven fields: label field, operation field, format field, tag field, operand field, comments field, and identification sequence field.

**Label Field (Columns 21-25)**

The label field represents the machine location of either data or instructions. The field may be left blank, may contain an asterisk in column 21, or may be filled with a symbolic address, left-justified in the field. Only data or instructions that are referred to elsewhere in the program need a label, although a label that is not further referred to is not an error.

A label can consist of up to five alphanumeric characters, beginning at the leftmost position of the label field. A label is always a symbol and must therefore conform to the rules for symbols (see Symbols). The example below shows the symbol ALPHA used as a label.



If the label field is left blank, it is ignored by the Assembler and has no effect on the assembled program. If column 21 contains an asterisk (\*), the entire statement is treated as comments and appears only in the listing. If the field contains a symbolic name (label), and the statement represents a standard machine language operation (Add, Store, etc.), the value assigned to the label is the address of the assembled instruction, which is equal to the value of the Location Assignment Counter (see Location Assignment Counter) at the time the statement is encountered by the Assembler. Values assigned to labels of the various assembler instructions are specified in the section entitled Assembler Instructions.

**Operation Field (Columns 27-30)**

Each machine instruction and assembler instruction has a unique mnemonic operation code associated with it. When a particular operation is to be represented, its mnemonic code must be punched, left-justified, in columns 27-30 of the source statement record.

Program \_\_\_\_\_

Date \_\_\_\_\_

Programmed by \_\_\_\_\_

Page No. \_\_\_\_\_ of \_\_\_\_\_

Label		Operation		F T		Operands & Remarks					Identification				
21	25	27	30	32	33	35	40	45	50	55	60	65	70	75	80

Figure 1. Coding Form

Format Field (Column 32)

The format field specifies the type of machine instruction being represented and, in the use of short (one-word) instructions, how the displacement field is to be handled. Any one of four entries is permitted: two for short instructions, one for a direct long (two-word) instruction, and one for an indirectly-addressed long instruction. For convenience, these formats are referred to by the character used to specify them, namely blank format, X format, L format, and I format.

Blank Format. A blank in the format field (column 32) signifies a short instruction except with some of the extended mnemonics provided with the Disk Monitor Assembler, in which case a blank format

field specifies a long instruction. Bit 5 of the assembled instruction is set to zero. A blank also indicates that any expression in the operand field be interpreted as the desired effective address for the statement.

During execution of certain short instructions, the effective address is the sum of the displacement (last 8 bits of the instruction word) and the contents of the Instruction Address Register (IAR). A blank format for such instructions causes the assembler to subtract the current value of the Location Assignment Counter from the expression in the operand field. Thus, when this result is added to the IAR during execution of the instruction, the correct effective address is obtained.

The effective address of short Store Index (STX) instructions is always obtained by adding the displacement to the IAR. The displacement of the Load



Index (LDX), Load Status (LDS), WAIT, all shift instructions, and all condition testing instructions is never added to the IAR. The effective address of all other short instructions is obtained by adding the displacement to the IAR, if the instructions are not indexed; that is, if column 33 is blank or zero.

The X format suppresses the automatic subtraction of the address counter from the displacement operand value when the instruction is moved. Therefore, the X format should be used for a short instruction which will have an effective address obtained by adding the displacement to the IAR. This requirement is not in conflict with the relocation process, because the process shifts the whole program, including instructions and reference data, to a core storage area different from that for which it was assembled. The relative distances between instructions and data remain the same, and the displacements remain correct.

In a relocatable assembly, the expression specifying an operand modified by the IAR must be relocatable so that the actual displacement is an absolute quantity (see Expressions). If this rule is not followed, a relocation error will be indicated. Also, since displacements must lie in the range  $-128_{10}$  to  $+127_{10}$ , the value of the displacement-specifying expression must not be more than  $127_{10}$  greater, nor more than  $128_{10}$  less than the address of the next location after the instruction in which it appears; otherwise, an addressing error will be indicated. An example illustrating the blank format is shown below:

Assume A = location  $1000_{10}$   
 B = location  $1050_{10}$

The value of the IAR will be  $1001_{10}$  when instruction A is executed. Therefore, the value computed by the assembler for the displacement will be  $49_{10}$ .

Label		Operation		F	T			
21	25	27	30	32	33	35	40	45
A		L, D,				B,		
		.						
		.						
		.						
		.						
		.						
B		D, C,				C, O, N, S, T,		

In the case of an instruction whose address is not modified by the IAR, the Assembler interprets the expression in the operand field as the desired contents of the displacement field, without modification. In this case, the operand specifying the displacement must be absolute and must be in the range  $-128_{10}$  to  $+127_{10}$ , or relocation and addressing errors result.

X Format. An X in the format field indicates to the Assembler that the related statement is to be assembled as a short instruction. It further indicates that any expression in the operand field is to be interpreted as the desired displacement value.

Consider the example illustrated in Figure 2; the purpose of this instruction sequence is to change the flow of a program by inserting a branch instruction in a location that previously contained a "no operation." If the branch instruction at BRCON were specified as MDX GO (i. e., blank format), the assembler would compute the displacement on the basis of the IAR value of 1101. (The IAR would have a value of 1101 if the BRCON instruction were executed where it was assembled.) However, the programmer, knowing the instruction will be executed at location SWTCH, computes the displacement himself and specifies the X format.

L Format. If column 32 contains the character L, it signifies a long (two-word) instruction with direct addressing. Bit 5 (F) of the assembled instruction is set to 1. The operand-field expression, which may be relocatable or absolute, is used to fill the second word (bits 16-31) of the assembled instruction. A second operand may be present, separated from the first operand by a comma (,). This operand may be used in one of two ways:

1. To specify symbolic condition codes for use with BSC, BSI and BOSC instructions.
2. To specify an expression that has a value in the range of  $-128$  to  $+127$  and is not relocatable.

This second operand yields bits to fill bit positions 8-15 of the assembled instruction.

I-Format. If column 32 contains the character I, it signifies an indirectly addressed long instruction. Bit 5 and bit 8 are set to 1. In all other respects an indirect instruction is treated exactly as a long direct instruction. If a displacement operand is specified, its high-order bit (bit 8) will always be a one, causing the displacement to be negative, because this bit is also the indirect flag bit.

Label	Operation	F	T	Operands & Remarks			Identification					
21 25	27 30	32 33	35	40	45	50	55	60	65	70	75	80
	•											
	•											
	•											
	•											
<i>S.W.T.C.H</i>	<i>N.O.P.</i>											
	•											
	•											
	<i>L.D.</i>					<i>B.R.C.O.N.</i>	<i>C.H.A.N.G.E</i>	<i>P.R.O.G.R.A.M</i>	<i>F.L.O.W</i>	<i>A.T.</i>	<i>S.W.T.C.H.</i>	
	<i>S.T.O.</i>					<i>S.W.T.C.H.</i>						
	•											
	•											
<i>B.R.C.O.N</i>	<i>M.D.X.</i>	<i>X</i>		<i>G.O.-</i>	<i>S.W.T.C.H.-1</i>							

Figure 2. Use of X Format

**Tag Field (Column 33)**

Column 33 is used to specify an index register if one is required. The code in column 33 is the index register number; i. e., 1=Index Register 1, 2=Index Register 2, and 3=Index Register 3. A zero or a blank indicates that no index register is to be used.

If no tag is specified in an LD<sub>X</sub>, MD<sub>X</sub>, or ST<sub>X</sub> instruction, the IAR is used. The example below shows an add instruction that addresses the core location whose address is zero plus the contents of Index Register 2.

Label	Operation	F	T	Operands & Remarks		
21 25	27 30	32 33	35	40	45	50
<i>S.U.M</i>	<i>A.</i>	<i>L2</i>		<i>0</i>		

**Operands and Remarks Field (Columns 35-71)**

The operand field is used to specify subfields in instructions and constants. The content of the operand

<i>S.H.O.R.T</i>	<i>S.T.O.</i>			<i>A.C.C.U.I.</i>
<i>L.O.N.G.</i>	<i>M.D.X.</i>	<i>L</i>		<i>A.C.C.U.I., 1, 0, 0, TWO-OPERAND, STATEMENT.</i>
<i>L.O.N.G.</i>	<i>S.T.O.</i>	<i>L</i>		<i>A.C.C.U.I., ONE-OPERAND, LONG, STATEMENT.</i>

Figure 3. One- and Two-Operand Statements

field for the various instruction formats are described under Format Field. Blanks must not appear within the operand(s) except as character values or in the EBC statements.

Some examples of one- and two-operand statements are shown in Figure 3.

**Remarks Field**

Remarks are for the convenience of the programmer. They permit lines or paragraphs of descriptive information about the program to be inserted in the program listing. Remarks appear only in the program listing; they have no effect on the assembled object program. Any valid characters (including blanks) can be used as remarks.

The Remarks field must appear to the right of the operand field and must be separated from it by at least one blank.

**Comments Field**

By placing an asterisk in column 21, the combined

statement fields from columns 22-72 may be used for comments. The identification-sequence field (columns 73-80) should not be used for comments.

If it is necessary to continue comments on additional lines, each line must have an asterisk in column 21, as illustrated in Figure 4.

#### Identification-Sequence Field (Columns 73-80)

The identification-sequence field may be used for program identification and statement-sequence numbers. It is limited to columns 73-80. The information in this field normally is punched in every statement card. The Assembler, however, does not check this field.

### STATEMENT WRITING

Symbolic language statements are accepted by the Assembler only if they conform to the rules of syntax presented in this section. Subsequent sections of this publication deal with the format and content of the specific types of assembler statements (machine instructions and assembler instructions). Instructions of both types are formed by using the basic elements described here. Many of the points introduced in this section are covered more extensively in subsequent sections.

#### Character Set

The following characters may be used in statements:

Monocase Alphabetics	A through Z, \$, #, @
Numerics	0 through 9
Special Characters	/*+ -= &   ^ < > ' . , ; ; ( ) % - ? (blank)

The codes that the assembler accepts for these characters are listed in Appendix A. Appendix A also lists additional codes which may be used in comments statements, as character values, and as alphameric constants. The + and & special characters may be used interchangeably as operators.

#### Symbols

Storage areas, instructions, and other elements may be given symbolic names for the purpose of referring to them in the program. The symbolic name is called a symbol. It can contain up to five characters.

While the first character of a symbol must be alphabetic, the remainder may be alphabetic, numeric, or any combination of the two. No embedded blanks or special characters may be used. Any violation of these rules is detected by the Assembler and indicated as an error in the program listing.

The following are valid symbols:

PUNCH	START	N
A2345	LOOP2	BC\$#@

\$. # and @ are monocase alphabets, not special characters (see Character Set), and as such can be used in the label field.

The following symbols are invalid, for the reasons noted:

256B	First character is not alphabetic
RECORDAREA2	More than 5 characters
END 1	Contains a blank

If a symbol is to be used as an operand, it must be defined in the program by using it as the label of a statement. Two types of label assignments are allowed. In machine-instruction statements and certain assembler statements, the label is assigned an address equal to the current value of the Location Assignment Counter. In the Equate Symbol statement (see Symbol Definition Statement), the label is assigned the value specified in the operand of the statement.

Symbol Table. For every program assembled, a table of the symbols in that program is created. This is the symbol table; each entry in the table records the value and relocation property of a symbol.

All symbols defined in the program are entered in the symbol table. Symbols that appear in the label field of assembler instructions that do not use labels (for example, ABS, END, ENT) are not placed in the symbol table.

General Restrictions on the Use of Symbols. The following restrictions are imposed on the use of symbols:

- A symbol may appear only once in a program as the label of a statement. If a symbol is used as a label more than once, only the first usage is recognized. Each subsequent usage of the symbol as a label is ignored and, in the card/paper tape system, is noted as an error in the program listing. In addition, any reference to

Label		Operation		F	T	Operands & Remarks							
21	25	27	30	32	33	35	40	45	50	55	60	65	70
*T,H,E,		S,T,E,R		S	K	I,N,C,O,L,2,I,M,A,K,E,S,T,H,I,S,A,C,O,M,M,E,N,T,S,L,I,N,E,							
*A,N,A		T,E,R,I		K		S,R,E,Q,U,I,R,E,D,F,O,R,E,A,C,H,L,I,N,E,O,F,C,O,M,M,E,N,T,S,							

Figure 4. Example of Comments Statement

such a symbol is noted as an error.

- The number of symbols that can be defined in a program is restricted by the amount of core storage available to the assembler (see IBM 1130 Card/Paper Tape Programming System Operator's Guide (C26-3629) or IBM 1130 Disk Monitor System, Version 2, Programming and Operator's Guide (Form C26-3717)).

### LOCATION ASSIGNMENT COUNTER

The Assembler maintains a counter to assign sequential storage addresses to program statements. This counter is called the Location Assignment Counter. It always indicates the next available address. As each machine instruction is processed, the counter is incremented by the number of words assigned to that instruction. Certain assembler instructions also cause the Location Assignment Counter to be set or incremented, whereas others do not affect it (see Assembler Instructions).

Location Assignment Counter Overflow. The maximum value of the Location Assignment Counter is 65535, a 16-bit value. If a program being assembled causes the counter to be incremented beyond 65535, the Assembler retains only the rightmost 16 bits in the counter and continues the assembly, checking for any other source program errors. No usable object program is produced. The user can, however, still obtain a listing of the entire source program.

### RELATIVE ADDRESSING

Once an instruction has been named by a symbol in the label field, it is possible for other instructions to refer to that instruction by using the same symbol. Moreover, it is possible to refer to instructions preceding or following the instruction named by indicating their positions relative to that instruction. This procedure is referred to as relative addressing. A relative address is, effectively, a type of expression (see Expressions).

For example, in the sequence

Label		Operation		F	T	Operands & Remarks							
21	25	27	30	32	33	35	40	45	50	55	60	65	70
S,T,A,R,T		A,				B,E,T,A							
		S,				S,T,O,R,E							
		S,T,O,		L		A,D,D,R,1							
A,L,I,S,T		A,		L		L,I,S,T							
		D,				L,O,C,2							

control can be transferred to the second instruction by either of the following instructions:

Label		Operation		F	T	Operands & Remarks							
21	25	27	30	32	33	35	40	45	50	55	60	65	70
		B,S,C,		L		S,T,A,R,T+1							
		B,S,C,		L		A,L,I,S,T-3							

By using relative addressing, it is also possible to refer to a particular word within a block of reserved storage. For example, the instruction

Label		Operation		F	T	Operands & Remarks							
21	25	27	30	32	33	35	40	45	50	55	60	65	70
B,E,T,A		B,S,S,				50							

reserves a block of 50 words, in which BETA is the address assigned to the first word in the block. The address BETA+1 then refers to the second word, BETA+2 to the third word, and BETA+n to the (nth+1) word.

Relative addressing can also be effected by using the current value of the Location Assignment Counter in an operand. In symbolic language this value is denoted by an asterisk (\*). (See The Asterisk Used as an Element.)

### SELF-DEFINING VALUES

A self-defining value is a machine value or a bit configuration.

Self-defining values can be used to specify such program elements as data, masks, addresses, and address increments. The type of representation selected (decimal, hexadecimal, or character) depends on what is being specified.

### Decimal Values

A machine decimal value is an absolute number from 0 to 65535. It is assembled as its binary equivalent. Some examples of decimal, self-defining values are

500	003
17	52324
7230	1

If a number larger than 65535 is specified in address arithmetic, the value is truncated modulo 65536; that is, only the low order 16 bits of the binary value are retained.

### Hexadecimal Values

A hexadecimal value is an unsigned hexadecimal number written as a sequence of digits. The digits must be preceded by a slash (/). The hexadecimal digits represent the 16 possible combinations of four bits.

Each hexadecimal digit is assembled as its four bit value. The hexadecimal digits and their bit patterns are as follows:

0 - 0000	4 - 0100	8 - 1000	C - 1100
1 - 0001	5 - 0101	9 - 1001	D - 1101
2 - 0010	6 - 0110	A - 1010	E - 1110
3 - 0011	7 - 0111	B - 1011	F - 1111

The following are examples of hexadecimal, self-defining values:

/FFFF	} equivalent
/AB12	
/379B	
/F2	
/00F2	

If more than four hexadecimal digits are specified in one sequence, only the four low-order digits are retained by the assembler. If less than four hexadecimal digits are specified, they are entered, right-justified.

A table for converting decimal values to hexadecimal values is provided in Appendix B.

### Character Values

A character value is a single character, preceded by a period. A character value may be a blank, any

combination of punches in a single card column, or a paper tape character that translates into the eight-bit IBM Extended BCD Interchange Code. Appendix A is a table of these combinations, their interchange codes and, where applicable, their printer graphics. A period used as a character value is represented as two periods in sequence, (i. e., ..).

Examples of character values are:

- . A
- . 1
- . 2
- . D
- . (blank)

The same value can frequently be represented by any one of the three types of self-defining values. For example, the decimal value 196 can be expressed in hexadecimal as /C4 and as a character, .D. The selection of a particular type of value is left to the programmer. Decimal values can be used for actual addresses and input/output unit numbers, hexadecimal values for masks, and character values for data.

## EXPRESSIONS

The term "expression" refers to symbols or self-defining values used as operands, either singly or in arithmetic combinations. Expressions are used to specify the various fields of machine instructions. They are also used as the operands of assembler-instruction statements.

An expression has three components: elements, terms, and operators.

### Elements

The smallest component of an expression is an element. An element is either a single symbol or a single self-defining value. The following are valid elements:

- TMP
- /1A6
- . B
- A
- \*
- 4

### The Asterisk Used As an Element

When used as an element the asterisk is relocatable and stands for the current value of the Location Assignment Counter for the instruction in which it appears (i. e., the rightmost word of the current instruction + 1). Thus, the asterisk as an element can have different values for different instructions.

Label		Operation				F	T			
21	25	27	30	32	33			35	40	45
		L,D,						A,B,C,		
S,U,M,		A,						D,E,F,		
		S,						D,A,T,A,		
		B,S,C,		L				S,U,M,	+	

The last instruction is a conditional branch to location SUM and can be written

Label		Operation				F	T			
21	25	27	30	32	33			35	40	45
		B,S,C,		L				*-4,	+	

Be sure the asterisk refers to the proper word when it is used with a long instruction or in an area where long instructions are present. In the previous example, the BSC instruction will become two machine language words after assembly. Therefore, during assembly of the BSC instruction, the Location Assignment Counter contains a value one greater than if the BSC were a short instruction.

### Terms

A term can consist of a single element, two elements separated by an asterisk (which denotes multiplication), or three elements each separated by an asterisk, etc. A term must begin with an element and end with an element, but is not permissible to write two elements in succession. The following are valid terms:

```
TMP * FUNC * TAXY
A * 4
X * Y * 5
6 * 4096
3
```

### Operators

An operator is a character that denotes an arithmetic function. The recognized operators are + or & (plus or ampersand), - (minus), and \* (asterisk), denoting addition, subtraction, and multiplication, respectively: An operator must be used between two terms. Two operators may not be used in succession.

There is no ambiguity between the use of the asterisk as an element and the use of the asterisk as an operator to denote multiplication, because the

position of the asterisk always makes clear what is meant. Thus, \*\*10 means "the value of the Location Assignment Counter multiplied by 10."

### Evaluation of Expressions

From a symbolically written operand, the evaluation procedure derives an integer value that can be used as (1) a displacement value in a short instruction, (2) an address in a long instruction, or (3) an absolute numeric quantity.

An expression is evaluated as follows:

1. Each element is replaced by its numeric value.
2. Each term is evaluated by performing the indicated multiplications from left to right, in the order in which they occur. In multiplication, the low-order 16 bits are retained.
3. The terms are combined from left to right, in the order in which they occur. If the result is negative, it is replaced by its 2's complement.

Grouping of terms, by parentheses or otherwise, is not permitted; however, this restriction can often be circumvented. For example, the product of 25 times the quantity B-C can be expressed as

$$25 * B - 25 * C$$

### Types of Expressions

In addition to evaluating expressions, the Assembler must decide whether the expression is absolute or relocatable. Without this information the Assembler would be unable to assign the proper relocation indicator bits for use during loading.

### Rules for Determining the Type of Expression

The rules by which the expression type is determined are:

- A symbol that is defined by means of the Location Assignment Counter is a relocatable element.
- Decimal and hexadecimal integers and character values are absolute elements.
- A relocatable element alone is a relocatable expression.
- A relocatable element, plus or minus an absolute element, is a relocatable expression.

- The difference of two relocatable elements is an absolute expression.
- A symbol that has been equated to an expression (by means of the EQU assembler instruction) assumes the same relocation property as that expression.

These rules are clarified by the following example:

Assume that a programmer wishes to incorporate a table into a relocatable program, and he knows that he may later wish to add or delete items without changing program references to the table. The first step is to assign symbols to the first (lowest-addressed) word in the table and to the location immediately after the last (highest-addressed) word of the table. These symbols could be BGTBL and ENTBL, respectively. Regardless of the number of items in the table or of the number of later additions or deletions, the number of words in the table is always equivalent to the value of the expression ENTBL-BGTBL. This illustrates the rule that the difference of two relocatable elements is an absolute expression.

Expanding this example, assume the programmer wishes to use a second table the same length as the first. The first (lowest addressed) word of the second table can be indicated by the symbol STBL. Then, the location following the last (highest-addressed) word of the second table can be indicated by the expression

$$STBL + ENTBL - BGTBL$$

This address is subject to relocation; hence, the expression is relocatable, following the rule that a relocatable element plus or minus an absolute element is a relocatable expression.

#### Procedure for Determining the Type of Expression

The following paragraphs describe the procedure for determining expression type (absolute or relocatable):

- Discard any term that contains only absolute elements.
- Examine each term of the expression. If any term contains more than one relocatable element, the expression will yield a relocation error.

- Replace each relocatable element by the symbol r, and replace each absolute element by its value. This yields a new expression which involves only numbers and the symbol r.
- Rewrite the expression in simplest form by evaluating it according to the address arithmetic rules given above in the section, Evaluation of Expressions.

If the result is an integer, the operand is absolute. If the result is r, the expression is relocatable. If the result contains r to any power other than one, or contains r with a coefficient other than one, the operand does not have a well-defined relocation property and will yield a relocation error. The following examples illustrate this procedure.

NOTE: When the terms absolute symbol and relocatable symbol are used in text, they mean symbols that refer to addresses.

Example 1: Consider the expression,

$$4+3*TRANS-2*FUNC+COUNT$$

where TRANS and FUNC are relocatable symbols, and COUNT is an absolute symbol. Discarding the terms involving only absolute elements leaves

$$3*TRANS-2*FUNC$$

This does not contain any illegal terms. Replacing each symbol by the letter r results in

$$3*r-2*r$$

Evaluating this produces r; therefore, the expression is relocatable.

Example 2: Consider the expression,

$$2*3*TRANS-FUNC$$

This reduces to

$$2*3*r-r$$

or

$$5r$$

This is neither  $r$  nor a number; therefore, the expression will cause a relocation error.

Example 3: Consider the expression,

$$A*2*R-A*A*R+5$$

where  $A$  is an absolute symbol, and  $R$  is a relocatable symbol. The expression is absolute if the value of  $A$  is zero or two and relocatable if the value of  $A$  is 1. If the value of  $A$  is anything else, a relocation error will result.

In the following examples,  $A$ ,  $B$ ,  $C$ , and  $D$  are relocatable symbols, and  $J$ ,  $K$ ,  $L$ ,  $M$ , and  $N$  are absolute symbols.

Relocatable expressions:

$A$	$1*A$
$A+J$	$250*A-249*B$
$A+B+C-D-*$	$100*A+50*B-75*C-74*D$

Absolute expressions:

$12345$	$0*A$
$A-B+C-D+5$	$500*A-400*B-100*C$

## Relocation Errors

If a source program contains an expression having in it one or more of the following, that expression is flagged as a relocation error.

- The negative (complement) of a relocatable element
- An absolute element minus a relocatable element
- The sum of two relocatable elements

In the following examples,  $A$ ,  $B$ ,  $C$ , and  $D$  are relocatable symbols, and  $J$ ,  $K$ ,  $L$ ,  $M$ , and  $N$  are absolute symbols.

$A+B$	$(+2r)$	$A*B$	$(r^2)$
$-A$	$(-1r)$	$2*A$	$(2r)$
$15-*$	$(-1r)$	$5*A-6*A$	$(-1r)$

$$A+J+M+N+B-C+D+L(+2r)$$

NOTE: In an absolute assembly headed by an ABS statement (described later), all symbols and asterisk values are defined as being absolute; therefore, no relocation errors are possible.



## MACHINE-INSTRUCTION STATEMENTS

All machine instructions can be represented symbolically as assembler language statements. There are two basic formats: short and long. However, within each basic format, further variations are possible.

The symbolic format of a machine instruction parallels, but does not duplicate, its actual format. A mnemonic operation code is written in the operation field, and one or more operands are written in the operand field. Comments can be appended to a machine-instruction statement as previously explained.

Any machine-instruction statement can be named by a symbol, which other assembler statements can use as an operand. The value of the symbol is the address of the leftmost word assigned to the assembled instruction.

### MNEMONICS

A list of all IBM 1130 machine language instructions and their associated mnemonics, including those mnemonics available for the monitor system only, is given in Table 1.

#### Condition-Testing Instructions (BSC, BOSC, BSI)

The machine instructions Branch or Skip on Condition (BSC), Branch Out or Skip on Condition (BOSC), and the long form of Branch and Store Instruction counter (BSI) use bits 10-15 of the displacement to test any combination of six conditions associated with the accumulator. When coding these instructions, the user does not use an expression to specify the displacement field, but, instead, writes a series of unique characters, each of which represents one bit of the condition-testing mask. These character symbols may be written in any combination; the bits they represent are combined by the assembler in a logical OR fashion. The symbols and their representations are:

Unique Character	Condition	Description	Bit Position Set to 1
O (Alpha)	Overflow	Skip or do not branch if Overflow indicator <u>off</u>	15
C	Carry	Skip or do not branch if Carry indicator <u>off</u>	14
E	Even	Skip or do not branch if bit 15 of Acc =0	13
+ or &	Plus	Skip or do not branch if bit 0 of the Acc =0, but not all bits of Acc =0	12
-	Minus	Skip or do not branch if bit 0 of Acc =1	11
Z	Zero	Skip or do not branch if all bits of Acc =0	10

Examples:

Operation	F	T			
27	30	32	33	35	
40					
BSC,			+	_____	Skip on plus condition
BSC,			+ -	_____	Skip on non-zero (plus or minus)
BSC,			Z -	_____	Skip on non-plus (zero or minus)
BSC,			C	_____	Skip if Carry indicator off
BSC,	L		EXIT, + -	_____	Branch to EXIT if not plus (zero or minus)
BSC,	L		EXIT, + -	_____	Branch to EXIT if zero (not plus or minus)
BSC,	L		EXIT	_____	Unconditional Branch to EXIT
BSC,	LI		0, Z+	_____	Branch to the contents of XR1 if minus (not zero or plus)
BSI,	L		SUBR, O	_____	Branch and Store instruction counter to SUBR if Overflow is on

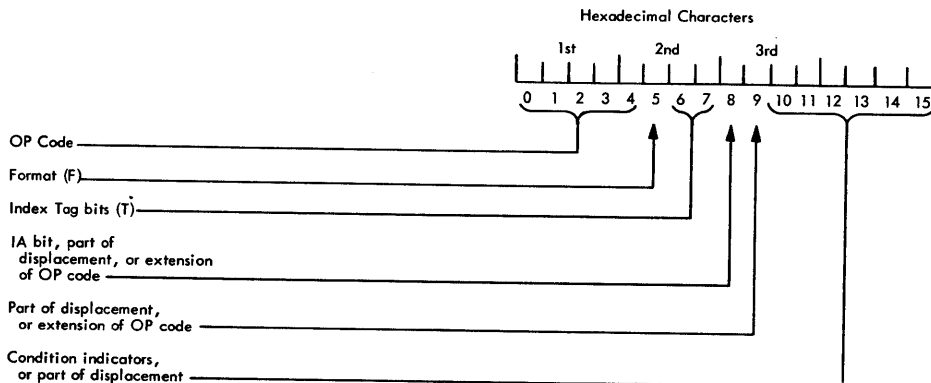
Table 1. Machine Instruction Mnemonics

Mnemonic	OP Code (Hexadecimal Representation) <sup>1</sup>	Instruction
<u>Load and Store</u>		
LD	C00	Load Accumulator
LDD	C80	Load Double
LDX	600	Load Index
LDS*	200	Load Status
STO	D00	Store Accumulator
STD	D80	Store Double
STX	680	Store Index
STS	280	Store Status
<u>Arithmetic</u>		
A	800	Add
AD	880	Add Double
S	900	Subtract
SD	980	Subtract Double
M	A00	Multiply
D	A80	Divide
AND	E00	And
OR	E80	Or
EOR	F00	Exclusive Or
MDM	† 5 740	Modify Memory
<u>Branch</u>		
B	† 4 700 or 4C0	Branch
BSI	400	Branch and Store Instruction Counter
BSC	480	Branch or Skip Conditionally
BP	† 6 4C30	Branch Accumulator Positive
BNP	† 6 4C03	Branch Accumulator Not Positive
BN	† 6 4C28	Branch Accumulator Negative
BNN	† 6 4C10	Branch Accumulator Not Negative
BZ	† 6 4C18	Branch Accumulator Zero
BNZ	† 6 4C20	Branch Accumulator Not Zero
BC	† 6 4C02	Branch on Carry
BO	† 6 4C01	Branch on Overflow
BOD	† 6 4C04	Branch Accumulator Odd
SKP*	† 480	Skip on Condition(s)
BOSC	2 484	Branch Out or Skip Conditionally
MDX	700	Modify Index and Skip
<u>Shift</u>		
SLA*	100	Shift Left Accumulator
SLT*	108	Shift Left Accumulator and Extension
SLC*	10C	Shift Left and Count Accumulator and Extension
SLCA*	104	Shift Left and Count Accumulator
SRA*	180	Shift Right Accumulator
SRT*	188	Shift Right Accumulator and Extension
RTE*	18C	Rotate Right
XCH*	† 3 18D	Exchange Accumulator and Extension
<u>Input/Output</u>		
XIO	080	Execute I/O
<u>Miscellaneous<sup>3</sup></u>		
NOP*	100	No Operation
WAIT*	300	Wait

\*Valid in short format only

†Not included in card/paper tape Assembler.

- The hexadecimal representation of the machine operation code is derived from the instruction format in the manner shown below. Bits 6 and 7 are assumed to be zeros because they do not enter into the makeup of any operation codes.
- Same as BSC with Bit 9 set to one.
- An operand should not be specified.
- When branch is short (Blank or X format), this operation code is assembled as an MDX (700). If the branch is long (L or I format), this operation code is assembled as a BSC with Bit 5 set to one (4C0).
- This instruction is automatically assembled as a long instruction (L is not required in the format field). Note that an attempt to use indirect addressing will result in a syntax error. Indexing is not permitted with this extended operation code.
- Extended conditional branch operation codes are assembled automatically as long instructions. (L is not required in the format field). Note that the proper condition code bits are preset, and further condition bits may not be specified following the operand.



## ADDITIONAL MONITOR SYSTEM MNEMONICS

Several new mnemonic operation codes which are equivalent to a Branch or Skip on Condition (BSC) may be used with the Monitor system. The operation code to be used for a specific job depends on the format and condition code required.

A new mnemonic MDM has been introduced that may be used in place of an unindexed MDX long. XCH may be used in place of RTE 16.

Examples of the additional Monitor System mnemonics are shown in Table 2. The mnemonics are listed below.

Skip on Condition (SKP). The condition codes (+, -, Z, E, O, and C) are specified as with a short BSC instruction. This instruction must not be indexed.

Branch Unconditionally (B). If the Format field contains an L or I, the BSC operation code is used with bit 5 set to one. Condition codes are not allowed after the address expression in the Operand field. If the Format field is left blank or contains an X, the MDX operation code is used, and the expression in the Operand field is used to form the displacement.

Branch Accumulator Positive (BP). Condition codes for accumulator zero (Z) and accumulator negative (-) are set to one.

Branch Accumulator Not Positive (BNP). Condition code for accumulator positive (+) is set to one.

Branch Accumulator Negative (BN). Condition codes for accumulator zero (Z) and accumulator positive (+) are set to one.

Branch Accumulator Not Negative (BNN). Condition code for accumulator negative (-) is set to one.

Branch Accumulator Zero (BZ). Condition codes for accumulator positive (+) and accumulator negative (-) are set to one.

Branch Accumulator Not Zero (BNZ). Condition code for accumulator zero (Z) is set to one.

Branch on Carry (BC). Condition code for Carry indicator off (C) is set to one.

Branch on Overflow (BO). Condition code for Overflow indicator off (O) is set to one.

Branch Accumulator Odd (BOD). Condition code for accumulator even (E) is set to one.

NOTE: Condition codes may not be used with any of the above instructions, except SKP, since the condition code is implicit in the extended mnemonic. The conditional branch instructions (all except SKP and B) are always assembled as long instructions; thus, the Format field need not contain an L, although the instruction is not classed as an error if L is specified. Indirect addressing may be specified.

Modify Memory (MDM). Contents of the location specified by the first operand is incremented or decremented by the value of the second operand. The second operand must be in the range -128 to +127.

NOTE: This instruction is always assembled as a long instruction; thus, the Format field need not contain an L, although the instruction is not classed as an error if L is specified. Indexing and indirect addressing must not be specified. If the operand becomes zero or changes sign, the next word in the program will be skipped.

Exchange Accumulator and Extension (XCH). Exchange is identical to a RTE of 16. No operand is specified with this instruction.

Table 2. Examples of New (Extended) Machine Instruction Mnemonics

New Instruction Statements					Equivalent Statements					Operations Performed		
27	30	32	33	35	40	27	30	32	33		35	40
S.K.P.				+		B.S.C.				+		Skip if accumulator is positive
S.K.P.				+ -		B.S.C.				+ -		Skip if accumulator is non-zero
S.K.P.				Z		B.S.C.				Z		Skip if accumulator is zero
S.K.P.				O		B.S.C.				O		Skip if Overflow indicator is off
S.K.P.				C		B.S.C.				C		Skip if Carry indicator is off
S.K.P.				+ - C		B.S.C.				+ - C		Skip if accumulator is non-zero or if Carry indicator is off
B.				EXIT		M.D.X.				EXIT		Branch unconditionally to EXIT, where EXIT must be within normal displacement range.
B.	L			ALPH		B.S.C.	L			ALPH		Branch unconditionally to ALPH
BZ				BETA		B.S.C.	L			BETA, + -		Branch to BETA if accumulator is zero
BN				BETA		B.S.C.	L			BETA, Z +		Branch to BETA if accumulator is negative
BNZ	I			BETA		B.S.C.	I			BETA, Z		Branch indirectly to BETA (i.e., the address specified by contents of BETA) if accumulator is non-zero
BN				RTNA		B.S.C.	L			RTNA, Z +		Branch to RTNA if accumulator is negative
BNN				RTNB		B.S.C.	L			RTNB, -		Branch to RTNB if accumulator is non-negative (zero or positive)
B.P.				SUB@		B.S.C.	L			SUB@, Z -		Branch to SUB@ if accumulator is positive
B.P.	I			SUB\$		B.S.C.	I			SUB\$, Z -		Branch indirectly to SUB\$ (i.e., the address specified by the contents of SUB\$) if accumulator is positive
BNP				SUB#		B.S.C.	L			SUB#, + -		Branch to SUB# if accumulator is non-positive (zero or negative)
BC				ENTR+1		B.S.C.	L			ENTR+1, C		Branch to ENTR+1 if Carry indicator is on
BC	I1			O		B.S.C.	I1			O, C		Branch indirectly to address specified by contents of index register 1 if Carry indicator is on
BO	Z			5		B.S.C.	L2			5, O		Branch to address specified by contents of index register 2 plus 5 if Overflow indicator is on
B.O.D.				\$AFE		B.S.C.	L			\$AFE, E		Branch to \$AFE if accumulator is odd
M.D.M.				SAVA, +5		M.D.X.	L			SAVA, +5		Increment contents of core location SAVA by 5
M.D.M.				/1D6A, 100		M.D.X.	L			/1D6A, 100		Increment contents of core location /1D6A by 100 decimal
M.D.M.				A, -12		M.D.X.	L			A, -12		Decrement contents of core location A by 12
X.C.H.						R.T.E.						Exchange the accumulator and extension (rotate right 16)

Just as machine instructions are requests to the computer to perform a sequence of operations during program execution, assembler instructions are requests to the Assembler to perform certain operations during the assembly. In contrast to machine-instruction statements, assembler-instruction statements do not always cause machine instructions to be included in the assembled program. Some, such as BSS and BES, generate no instructions but do cause storage areas to be set aside for constants and other data. Others (e. g., EQU) are effective only during the assembly; they may or may not generate something in the assembled program. If nothing is generated, the Location Assignment Counter is not affected.

The following is a list of all assembler statements permitted by the IBM 1130 Card/Paper Tape Assembler. These statements are also valid for the Monitor Assembler. Additional statements are provided for the Monitor Assembler and are listed in the section Monitor Assembler Statements.

**Program Control**

- ABS - Absolute Assembly
- LIBR - Transfer Vector Subroutine
- SPR - Standard Precision
- EPR - Extended Precision
- ORG - Define Origin
- END - End of Source Program

**Data Definition**

- DC - Define Constant
- DEC - Decimal Data
- XFLC - Extended Floating Constant
- EBC - Extended Binary Coded Information

**Storage Allocation**

- BSS - Block Started by Symbol
- BES - Block Ended by Symbol

**Symbol Definition**

- EQU - Equate Symbol

**Program Linking**

- ENT - Define Subroutine Entry Point
- ISS - Define Interrupt Service Entry Point
- ILS - Define Interrupt Level Subroutine
- CALL - Call Subroutine (2-word call)
- LIBF - Call Subroutine (1-word call)

**PROGRAM CONTROL STATEMENTS**

Program control statements are used to set the Location Assignment Counter to a specific value, to define the end of a source program, or to specify whether a particular program is to be assembled as absolute or relocatable. None of these assembler statements generate machine-language instructions or constants in the object program.

ABS — Assemble Absolute

An ABS statement is used to specify that a main program is to be assembled as an absolute program. An absolute program is one in which the core locations used at execute time are the same as those specified by the programmer in the source program. The ABS statement is punched as shown below and is then used as the first statement of a source program.

21	25	27	30	32	33	35	40	45
		A,B,S						

If the first (non-comment) statement of a source program is not an ABS statement, the program will be assembled as relocatable. In an absolute assembly headed by an ABS statement, all symbols and asterisk values are defined as absolute quantities; therefore, no relocation errors are possible. The significance of relocatable and absolute assemblies is explained in the following paragraphs.

**Relocatable Assembly**

Some programs assembled by the IBM 1130 Assembler are absolute; that is, the locations of assembled instructions are known during the assembly and the location on the listing is the actual location where a particular word is loaded. However, subroutines used by an absolute program must be in such a form that they may be loaded at various locations; otherwise, it would be necessary for the user to reassemble the subroutines each time he assembled a main program that required them. Therefore, all subroutines must be and main programs may be assembled relocatable.

Every relocatable program or subroutine produced by the IBM 1130 Assembler is assembled as though it begins at location zero. Since a job to be executed may contain several subroutines, it is obvious that they cannot all be loaded into locations starting with location zero. In fact, no relocatable program is ever loaded at location zero; instead, each program is relocated. The relocatable main program is loaded into the first available location. Subroutines are then loaded into successively higher locations of core storage, each beginning with the

next even location after the last core storage location used by the preceding subroutine. When a particular program has been loaded, the address of the first word is called the load address for that program.

Thus, the address in core storage actually occupied by an instruction of the program is the address assigned to that instruction during assembly, plus the load address of that program. To keep the program self-consistent, the load address must be added to the address of many (but not all) 2-word instructions, and those constants whose values are relocatable.

This process of conditionally adding the load address is performed by the loading program before execution and is called relocation. In relocating instructions, the loading program is guided by relocation indicator bits which are a part of the object program.

### Absolute Assembly

The programmer uses the ORG assembler statement in his source program to specify the locations into which the object program resulting from an absolute assembly is loaded. Subroutines are loaded into successively higher even-core locations following the end of the main program.

Only main programs may be assembled absolute; subroutines must be assembled relocatable.

### LIBR — Transfer Vector Subroutine

An LIBR statement is used as the first statement of a subroutine to specify that the subroutine is to be called by LIBF statements only (see Program-Linking Statements). The absence of an LIBR statement specifies that the subroutine is to be called by CALL statements only. LIBR statements are for subroutines only, as ABS statements are for main programs only. An LIBR statement needs no operands.

### SPR — Standard Precision, EPR — Extended Precision

The SPR or EPR statement specifies that the program (main or subroutine) in which it appears uses standard precision or extended precision, respectively, for arithmetic operations. If these statements are included in the user's programs, the loader ensures that main programs and subroutines always match with regard to precision. Their use is optional, however.

If used, the SPR or EPR statement must follow the ABS or LIBR statement. If no ABS or LIBR statement is used, the SPR or EPR statement is the first statement in the program.

### ORG — Define Origin

This assembler instruction is used to set the Location Assignment Counter (i.e., the next location to be assigned) to any desired value. In this way the programmer is able to control the assignment of storage to instructions, constants, and data. If a Define Origin statement is not the first entry in an absolute source program, the processor begins the assignment of storage at a location compatible with the size of the applicable loader (Card/Paper Tape Assembler) or the version of disk I/O required (Disk Monitor Assembler). A typical Define Origin statement is shown below.

Label		Operation		F	T				
21	25	27	30	32	33	35	40	45	
		ORG,				3,0,0,0,			

The label, if used, is assigned a value equal to the value of the Location Assignment Counter at the time the statement is encountered in the source program. (This assignment is made before the counter is modified.) If any symbols are used in the expression, they must have been previously defined. In a relocatable assembly, an absolute expression in the operand field is considered a relocation error and the statement is ignored.

Some examples of Define Origin statements are given below:

Label		Operation		F	T					Op
21	25	27	30	32	33	35	40	45	50	
		ORG,				XYZ,				
S,T,A,R,T		ORG,				XYZ+50,				
S,T,A,R,T		ORG,				*+50, LOC, CTR,+50,				

If the label XYZ has been previously defined as 1000<sub>10</sub> the first entry directs the assembler to begin the assignment of succeeding entries at location 1000. The second entry directs the Assembler to begin the assignment of succeeding entries 50 core locations beyond the location that has been assigned to the symbol XYZ. The third entry directs the Assembler to begin the assignment of succeeding entries at the

address specified by the current address of the Location Assignment Counter plus 50.

### END — End of Source Program

An END statement is the last statement of a source program; it indicates to the assembler that all statements of the source program have been processed. An END statement is also used to define the execution address of the main program. To do this, the END statement requires an operand that represents the starting address of the program. At the completion of loading, execution begins at the address specified by the operand. For subroutines, all entry points are specified by ENT statements (described later); therefore, the operand of the END statement for a subroutine is blank.

The following statements illustrate both types of END statements.

Label	Operation	F	T	Operands
21	25	27	30	32 33 35 40 45 50
	END			END OF PROGRAM
	END			G.O. BRANCH TO G.O.

### DATA DEFINITION STATEMENTS

Data Definition statements are used to enter data constants into storage. The statements can be named by symbols so that other program statements can refer to the fields generated. Any type of data definition statement can be used in standard or extended precision program.

### DC — Define Constant

The Define Constant statement is for generating constant data in main storage. Data can be specified as characters, hexadecimal numbers, decimal numbers, storage addresses, or any valid expression. One 16-bit word is generated for each DC statement. The format of this statement is shown below:

Label	Operation	F	T	Operands
21	25	27	30	32 33 35 40 45
LABEL	DC			AN EXPRESSION

If a label is used, the address assigned to it is the location of the generated data word and is equal to the current value of the Location Assignment Counter. Some examples of DC statements follow:

Label	Operation	F	T	Operands
21	25	27	30	32 33 35 40 45 50
H,EX	DC			/FFFF, H,EX, CONST
D,EC	DC			-3,8,5, D,EC, INT,GER
A,L,P,H,A	DC			.B, C,H,A,R, C,O,N,S,T
A,D,D,R,S	DC			A,L,P,H,A+5, A,D,D,R, C,O,N

### DEC — Decimal Data

The Decimal Data statement is used to enter binary data, expressed in decimal form, into a program. One DEC statement generates two 16-bit words of binary information. The format of the DEC statement is as follows:

Label	Operation	F	T	Operands & Remarks
21	25	27	30	32 33 35 40 45 50
LABEL	DEC			Decimal Data Item

If a label is used, its value is equal to the current value of the Location Assignment Counter if the current value is even; if the current value is odd, the label will be equal to the current value plus one. The label is assigned to the leftmost word of the generated constant. The types of data permitted in the operand field are described in the paragraphs entitled Decimal Data Items. An example of a DEC statement follows:

Label	Operation	F	T	Operands
21	25	27	30	32 33 35 40 45
D,A,T,A	DEC			7,19

If the value of the Location Assignment Counter is 1000 when the DEC statement is encountered, the two words in storage look like this:

Location	Contents in Hexadecimal Form
01000	0000
01001	0013

### Decimal Data Items

A decimal data item is used to specify, in decimal form, two or three words of data to be converted into binary form. Decimal data items are used in the

operand field of DEC assembler statements. Three types of decimal-data items are permitted: decimal integers, real numbers, and fixed-point numbers. A real decimal-data item can also be used as the operand of an XFLC statement that generates a 3-word constant.

Decimal Integers. A decimal integer is composed of a series of numeric digits with or without a preceding plus or minus sign. The allowable range of decimal integers is  $-(2^{31}-1)$  to  $2^{31}-1$ .

Examples

<u>Decimal Integer</u>	<u>Stored As</u>
50	00000032 <sub>16</sub>
1535	000005FF <sub>16</sub>
-3729	FFFFF16F <sub>16</sub>
	(2's complement)

Real Numbers. A real number has two components: a mantissa and an exponent.

- Mantissa — The mantissa is a signed or unsigned decimal number, which can be written with or without a decimal point. The decimal point can appear at the beginning, at the end, or within the decimal number. If the exponent (see below) is present, the decimal point can be omitted, in which case it is assumed to be located at the right-hand end of the decimal number.
- Exponent — The exponent consists of the letter E, followed by a signed or unsigned decimal integer. The exponent part can be omitted if the mantissa contains a decimal point. If used, it must follow the mantissa.

A real number is converted to a normalized, real binary number. The exponent part, if present, specifies a power of ten by which the mantissa is multiplied during conversion. For example, all of the following real numbers are equivalent and will be converted to the same real binary number.

4.500  
 45.00E-1  
 4500E-3  
 .4500E1

In standard precision, the above real numbers are converted and stored in two consecutive storage locations as follows:

Word 1  
 4800

Word 2  
 0083

The DEC assembler instruction stores real numbers in the standard precision real number format described in the manual, IBM 1130 Subroutine Library (Form C26-5929).

Fixed Point Numbers. A fixed-point number can have up to three components: a mantissa, an exponent, and a binary-point identifier.

- Mantissa — The mantissa is the same as described for real numbers.
- Exponent — The exponent is the same as described for real numbers.
- Binary-Point Identifier — This identifier consists of the letter B, followed by a signed or unsigned decimal integer. The binary-point identifier must be present in a fixed-point number and must come after the mantissa. If the number has an exponent, the binary point identifier may precede or follow the exponent.

A fixed-point number is converted to a fixed-point binary number that contains an understood binary point. The purpose of the binary-point identifier of the number is to specify the location of this understood binary point within the word. The number that follows the letter B specifies the number of binary places in the word to the left of the binary point (that is, the number of integral places in the word). The sign bit is not counted. Thus, a binary-point identifier of zero specifies a 31-bit binary fraction. B2 specifies two integral places and 29 fractional places. B31 specifies a binary integer. B-2 specifies a binary point located two places to the left of the leftmost bit of the word; that is, the word would contain the low-order 31 bits of binary fraction. As with real numbers, the exponent, if present, specifies a power of ten by which the mantissa is multiplied during conversion.

A fixed-point number preceded by a minus sign is stored in 2's complement form.

The following fixed-point numbers all specify the same configuration of bits, but not all of them specify the same location for the understood binary point:

22.5B5  
 11.25B4  
 1125B4E-2



1125E-2B4  
9B7E1

All of the above fixed-point numbers are converted to the same binary configuration, whose hexadecimal representation is:

Word 1
5A00

Word 2
0000

**XFLC - Extended Real Constant**

The XFLC assembler instruction is used to introduce into a program an extended precision real constant, expressed in three consecutive data words. When assembled, this instruction produces a format identical to the extended range real format described in the manual, IBM 1130 Subroutine Library (Form C26-5929).

The format of the XFLC instruction is shown below:

Label	Operation	F	T	Operands & Remarks
21	25	27	30	32 33 35 40 45 50
L.A.B.E.L	X.F.L.C			R.E.A.L. N.U.M.B.E.R.

The label is optional; if it is used, it is assigned to the location of the leftmost word generated.

Some examples of the XFLC instruction are shown below:

Label	Operation	F	T	Operands & Remarks
21	25	27	30	32 33 35 40 45 50
	X.F.L.C			0.53125
R.E.A.L.	X.F.L.C			-0.53125
	X.F.L.C			5.12E2

The data (in hexadecimal form) generated by each of these examples is

1.    Word 1    Word 2    Word 3  
      0080    4400    0000
2.    Word 1    Word 2    Word 3  
      0080    BC00    0000
3.    Word 1    Word 2    Word 3  
      008A    4000    0000

**EBC - Extended Binary Coded Information**

The EBC statement is used to generate data words, each consisting of two 8-bit characters in the Extended BCD Interchange Code (see Appendix A). Up to 18 sixteen-bit words can be generated with one EBC statement. The format of the statement is shown below:

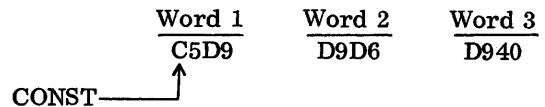
Label	Operation	F	T	Operands & Remarks
21	25	27	30	32 33 35 40 45
L.A.B.I.	E.B.C.			.A.L.P.H.A. .D.A.T.A.

If a label is present, it is assigned to the location of the leftmost word generated. The operand field contains the alphanumeric data to be represented in storage. This data must begin and end with a period. The data can be any valid character in the Extended BCD Interchange Code, including the period.

**Examples**

Label	Operation	F	T	Operands & Remarks
21	25	27	30	32 33 35 40 45
C.O.N.S.T	E.B.C.			.E.R.R.O.R.
A.L.P.H.A	E.B.C.			.C.O.N.S.T.A.N.T.

The first example generates three words of data, with the location of the label CONST assigned to the leftmost location of the first word generated.



Note that if the constant has an odd number of characters, as in the above example, the last word of data ends with the 8-bit equivalent of blank.

The second example generates four words of data:

Word 1	Word 2	Word 3	Word 4
C3D6	D5E2	E3C1	D5E3

NOTE: A period may not appear in the remarks field of an EBC instruction.

## STORAGE ALLOCATION STATEMENTS

Storage allocation statements are used to reserve blocks of storage for data or work areas. Two such statements are available with the IBM 1130 Assembler: Block Started by Symbol and Block Ended by Symbol.

### BSS — Block Started by Symbol

The BSS assembler instruction is used to reserve an area of core storage, within a program, for data storage or for working space. The format of the BSS instruction follows:

Label	Operation	F	T	Operands & Res
21 25	27 30	32 33	35	40 45 50
LABEL	BSS			Absolute Expression

The expression specifies the number of words to be reserved; the label, if specified, refers to the leftmost word reserved. The location of the block of storage within the object program is determined by the location of the BSS statement within the source program.

If the character E is punched in column 32, the assembler assigns the leftmost word of the reserved location to the next available even location. If a blank or any character other than E appears in column 32, the assembler assigns the leftmost word of the reserved area to the next available location regardless of whether that location is even or odd. This feature is useful when defining areas for use with double precision instructions.

A BSS statement with an E format and an operand value of zero causes the Location Assignment Counter to be made even (if necessary) before the next instruction is assembled.

A BSS instruction causes an area to be reserved, not cleared; therefore, it should not be assumed that an area reserved by a BSS instruction contains zeros.

Any symbols in the operand field of a BSS assembler instruction must have been previously defined. The expression in the operand field must be an absolute expression.

In the following example, the symbol AREA is equivalent to 3000; the next location assigned is 3028.

Label	Operation	F	T	Operands & Res
21 25	27 30	32 33	35	40 45
	ORG			3000
AREA	BSS			28

### BES — Block Ended by Symbol

The BES instruction is identical to the BSS instruction except that the address assigned to the label is the rightmost word in the area plus 1, i.e., the next location available for assignment.

In the previous example, the symbol AREA is equivalent to 3028.

## SYMBOL DEFINITION STATEMENT

One symbol definition statement (EQU) is available in the IBM 1130 Assembler language.

### EQU — Equate Symbol

The EQU statement is used to assign to a symbol a value other than the value of the Location Assignment Counter at the time the symbol is encountered. The format of the EQU statement is

Label	Operation	F	T	Operands & Res
21 25	27 30	32 33	35	40 45 50
SYMBOL	EQU			An Expression

The symbol in the label field is made equivalent to the value of the expression. The expression may be absolute or relocatable. All symbols appearing in this expression must have appeared as a label in a previous statement. If an asterisk (\*) is used as the expression, the value assigned to it is the next location to be assigned by the assembler.

### Examples

Label	Operation	F	T	Operands & Res
21 25	27 30	32 33	35	40 45
NAME	EQU			26
LOOP	EQU			NAME+1

In the first example, the symbol NAME is assigned a value of 26. In the second example, the symbol LOOP is assigned a value of 27.

Word 30 of the header record can be set for identification purposes as shown below. Word 30 is not used by any of the 1130 programs.

### LINKING STATEMENTS

Linking statements are used to establish communication between a main program and its subroutines or between a program and the Monitor system.

#### ENT - Define Subroutine Entry Point

The ENT statement should be used to define the entry point(s) in all subroutines except ISS and ILS. Up to fourteen entry points (ten with the Card/Paper Tape Assembler) may be defined for each subroutine (this would require an equal amount of ENT statements). The format of the ENT statement is shown below.

Label	Operation	F	T
21 25	27 30	32 33	35 40
_____	ENT		NAME
_____			
_____			

NAME is a symbol that identifies an entry point for the associated subroutine. This symbol must be relocatable. All ENT statements for a given subroutine must be together and must precede all statements except LIBR, SPR, EPR, and comments statements. ENT, ISS, or ILS statements (see below) may not be used in the same subroutine.

#### ISS - Define Interrupt Service Entry Point

IBM provides interrupt service subroutines (ISS) for all devices; however, the user is given the option of replacing or adding to these subroutines with his own. The ISS statement is used to define an entry point in an interrupt service subroutine and to establish interrupt linkages to the subroutine during loading. Only one entry point may be defined for each subroutine. The format of the ISS statement is shown below.

Label	Operation	F	T
21 25	27 30	32 33	35 40 45
_____	ISS	NN	NAME L
_____			
_____			

<u>Label</u>	<u>ISS Header Word 30</u>
blank	blank
1130	1
1800	2

NAME is as described for the ENT statement and NN (the ISS number) is a decimal number from 01 to 20 used during loading to establish the linkage from the appropriate point in the corresponding ILS. The numbers and associated devices used in the subroutines provided by IBM are listed below.

#### Card/Paper Tape System.

<u>Number*</u>	<u>Device or Function</u>
01	1442 Card Read Punch
02	Input Keyboard/Console Printer
03	1134 Paper Tape Reader; 1055 Paper Tape Punch
05	Single Disk Storage
06	1132 Printer
07	1627 Plotter

\*Numbers 08 through 20 are assignable by the user.

#### Monitor System.

<u>Number*</u>	<u>Device or Function</u>
01	1442 Card Read Punch; 1442 Card Punch
02	Input Keyboard/Console Printer
03	1134 Paper Tape Reader; 1055 Paper Tape Punch
04	2501 Card Reader
05	Single Disk Storage; 2310 Disk Storage
06	1132 Printer
07	1627 Plotter
08	Synchronous Communications Adaptor
09	1403 Printer
10	1231 Optical Mark Page Reader

\*Numbers 11 through 20 are assignable by the user.

NOTE: User-assigned ISS numbers should start at twenty and proceed backwards in order to avoid conflict with IBM-assigned ISS numbers.

L is a one-digit number required by the Card/Paper Tape Assembler to indicate the interrupt level(s) associated with the subroutine. The level numbers (0-5) can be listed in any order in columns 45, 50, 55, 60, 65, and 70 with the first appearing in 45, the second in 50, etc.

L is not used with the monitor system. Instead, LEVEL control cards are used with the subroutine being assembled, one card per interrupt level required (see Assembler Control Records in the publication IBM 1130 Disk Monitor System, Version 2, Programming and Operator's Guide (Form C26-3717)).

An ISS statement must precede all statements except LIBR, SPR, EPR and comments statements.

Procedures for writing ISSs are provided in the publications IBM 1130 Subroutine Library (Form C26-5929) and IBM 1130 Disk Monitor System, Version 2, Programming and Operator's Guide (Form C26-3717).

### ILS - Define Interrupt Level Subroutine

IBM provides interrupt level subroutines for the various I/O devices and their associated interrupt levels; however, the user may replace or add to these subroutines with his own. The ILS statement is used to define an interrupt level subroutine and to associate the subroutine with a specific interrupt level. The format of the ILS statement is shown below.

Label	Operation	F	T
21 25	27 30	32 33	35
	I,L,S		NN

NN is the interrupt level number (00-05) associated with the interrupt level subroutine and is used during loading. The devices associated with each interrupt level are shown below:

<u>Interrupt Level</u>	<u>Device(s)</u>
00	1442 Card Read Punch (1442 Card Punch)
01	1132 Printer (Synchronous Communications Adaptor)
02	Single Disk Storage (2310 Disk Storage)

### Interrupt Level

### Device(s)

03	1627 Plotter
04	Keyboard/Console Printer, 1442 Card Read Punch, 1134 Paper Tape Reader, 1055 Paper Tape Punch (2501 Card Reader, 1403 Printer, 1231 Optical Mark Page Reader)
05	PROGRAM STOP Key or Interrupt Run Mode.

NOTES: 1. The devices listed in parentheses are used with the Monitor system only.

2. An ILS statement must precede all statements except SPR, EPR, and comments statements.

Procedures for writing interrupt level subroutines are provided in the publications, IBM 1130 Subroutine Library (Form C26-5929) and IBM 1130 Disk Monitor, Version 2, Programming and Operator's Guide (Form C26-3717).

### CALL - Call Direct Reference Subroutine

A CALL statement is used to call some of the subroutines in the IBM Subroutine Library or any user-written subroutine written for the CALL statement. During execution, this type of call takes the form of a long (two-word) BSI (direct for card/paper tape system, indirect for Monitor system), to the entry point named in the CALL and the corresponding ENT or ISS statement.

When BSI is executed, the location of the first word following it is placed in the entry point location, and control is transferred to the first word following the entry point. The format of the CALL statement is:

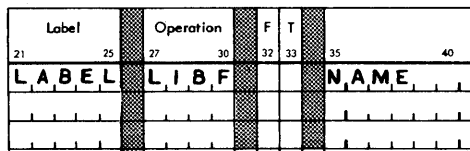
Label	Operation	F	T
21 25	27 30	32 33	35 40
LABEL	CALL		NAME

If used, the label is assigned to the current value of the Location Assignment Counter, which is the same as the leftmost word of the generated

BSI instruction. The name of the called subroutine is assembled into the object program, together with a unique code identifying the CALL. This code is used during loading to generate the BSI to this subroutine.

LIBF - Call TV (Transfer Vector) Reference Subroutine

An LIBF statement is used to call any of the subroutines in the Subroutine Library (or any user-written subroutine) written to utilize the Transfer Vector (see the following section). The format of the LIBF statement is:



If used, the label is assigned to the current value of the Location Assignment Counter when the LIBF statement is encountered. The name of the called subroutine is assembled into the object program, together with a unique code identifying the call as an LIBF call. This code is used during loading to generate the linkage to the subroutine. During execution, the TV subroutine uses Index Register 3. Therefore, if Index Register 3 is used by any other instruction in the user's program, it must be saved and restored before it is needed by any TV subroutine calls.

**LIBF Subroutine Transfer Vector**

To fully understand the use of the LIBF statement, the user should be familiar with the makeup of the transfer vector, which allows main programs to communicate with relocatable subroutines (and relocatable subroutines to communicate with each other) without knowing where in core storage the subroutines are loaded. The Transfer Vector consists of three 16-bit words for each subroutine entry point referred to by an LIBF statement. The contents of the three words vary as the subroutine goes through the three phases of being called, loaded, and executed. The following paragraphs describe these three phases, and illustrate the contents of the transfer vector for each phase.

Recognizing the Subroutine Call. All subroutines that utilize the Transfer Vector are called via LIBF statements. These statements take the following general form:

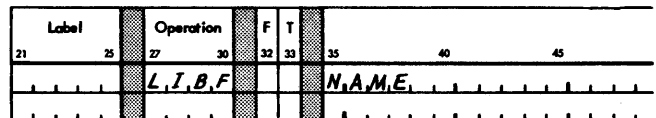
```
LIBF      NAME
DC        Parameter
DC        Parameter
etc.
```

When an LIBF call is recognized during the loading of an object program, the loader begins to build the transfer vector by saving the name of the called subroutine.

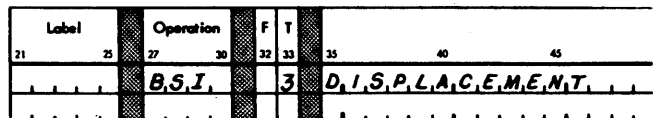


Subsequent LIBF statements produce additional records for the Transfer Vector, each containing a unique subroutine name. Calls to a subroutine previously listed in the transfer vector do not produce a new record. Ultimately each causes a short, indexed BSI instruction pointing to the first word of the associated Transfer Vector entry. This instruction, generated during loading, uses Index Register 3 and a computed displacement to refer to the proper Transfer Vector entry.

Original Statement



Modified Statement



When this BSI instruction is encountered during execution of the main program, it causes a branch to the associated Transfer Vector entry and from there to the entry point of the subroutine (see the following section, Loading the Subroutine). A BSI statement is generated for each LIBF statement encountered.

NOTE: Index Register 3 is reserved for LIBF subroutine calls. Therefore, if any instructions are to use Index Register 3, it should be restored prior to any LIBF subroutine call.

## MONITOR ASSEMBLER STATEMENTS

In addition to the basic assembler statements, the IBM 1130 Monitor Assembler is provided with the following capabilities.

### Disk Data Organization

- DSA - Define Sector Address
- FILE - Define Disk File

### Data Definition

- DMES - Define Message
- DN - Define Name

### Linking

- LINK - Load and Execute Another Program
- EXIT - Return Control to Supervisor
- DUMP - Dump and Terminate
- PDMP - Dump and Continue

### List Control

- HDNG - Print Heading on Each Page
- LIST - List Segments of Programs
- SPAC - Space Listing
- EJCT - Start New Page

## DISK DATA ORGANIZATION STATEMENTS

### DSA - Define Sector Address

The DSA statement allows the programmer to refer symbolically to a disk-stored data file or program stored in Disk Core Image format (DCI) without knowing the specific disk location of the data or program. The disk location of data files and programs can vary on disk because of deletions, but the DSA statement allows easy reference through the use of the symbolic name of the data file or program.

The format of the DSA statement is:

Label	Operation	F	T	Operands & Ren
21	25	27	30	32 33 35 40 45 50
LABEL	DSA			NAME

The label is defined as the current value of the Location Assignment Counter when the DSA statement is encountered. The symbol in the operand field must be the name of a data file or DCI program that is on disk both when the assembly is made and during execution.

The following statements illustrate the use of the DSA statement to read one sector of data. For a description of the disk calling sequences, see the publication IBM 1130 Subroutine Library (Form C26-5929).

Label	Operation	F	T	Operands & Ren
21	25	27	30	32 33 35 40 45 50
	.			
	LIBF			DISK1
	DC			1,0,0
	DC			I,O,A,R
	DC			E,R,R,O,R
	.			
I,O,A,R	DSA			DATA
	B,S,S			3,1,9
	.			
	.			

The Assembler reserves three words in the object program for each DSA statement. These words are filled in by the Core Load Builder. For a data file they will contain:

- Word 1 - Length (in words)
- Word 2 - Sector Address, including the drive code
- Word 3 - Sector count of the file

For a program they will contain:

- Word 1 - Length (in words)
- Word 2 - Sector Address, including the drive code
- Word 3 - Execution Address of the Program

If the area corresponding to the DSA statement is used as the I/O area for a disk read operation, the execution address of the program must be saved prior to the disk call to bring in the program. (The contents of the third word are destroyed by the incoming data).

The following statements illustrate the use of the DSA statement to supply the disk address of a one-sector program.

Label	Operation	F	T	Operands & Ref
21	25	27	30	32 33 35 40 45 50
	.			
	.			
	L.D.			I.O.A.R.+2
	S.T.O.			BRNCH.+1
READ	L.I.B.F			DISK1
	D.C.			/1000
	D.C.			I.O.A.R
	D.C.			ERROR
CALL	L.I.B.F			DISK1
	D.C.			/0000
	D.C.			I.O.A.R
	M.D.X.			CALL
BRNCH	B.S.C.	L		0
	.			
	.			
	.			
I.O.A.R	DSA			PRGRM
	B.S.S.			319
	.			
	.			

The following statements can be added to the previously shown program call to call a second program and have it loaded to the same area as the first.

Label	Operation	F	T	Operands & Ref
21	25	27	30	32 33 35 40 45 50
	L.D.			A.D.R.2
	S.T.O.			I.O.A.R
	L.D.			A.D.R.2+1
	S.T.O.			I.O.A.R+1
	L.D.			A.D.R.2+2
	S.T.O.			BRNCH.+1
	M.D.X.			READ
A.D.R.2	DSA			PRGRM2
	.			
	.			

The execution address of the second program can be different from the first, but the programs must be executable from the same locations. This requires a certain amount of planning before assembling the "overlay" programs.

### Programming Considerations

The following considerations must be observed by the user who wishes to use the DSA statement to supply the disk address for programs.

- The called programs must be in DCI format.
- If the calling program is converted to DCI format, the data for the DSA statement is filled in during the core image conversion and will be fixed for all subsequent executions. Thus, if the referenced program or data files are subsequently moved, incorrect results will occur. Data files referenced by a Core Image program should be stored in the Fixed area.
- Any loading functions, such as the setting of Index Register 3, will have to be supplied by the calling program.

### FILE - Define Disk File

The FILE statement specifies to the Assembler the file identification, the number of file records in a file, and the size of each record in a disk data file that will be used with a particular mainline and its associated subprograms. The Assembler FILE statement allows the Assembler language user to define files so that they are similar to FORTRAN defined files.

As a core load is constructed by the Core Load Builder, the defined files are equated to data files already assigned in the User/Fixed Area or to files in Working Storage.

The FILE statement must not appear in a subprogram; it is permitted only in a relocatable mainline program. Therefore, all subprograms used by the mainline must use the defined files of the mainline. The format of the FILE statement is as follows:

Label	Operation	F	T	Operands & Ref
21	25	27	30	32 33 35 40 45 50
l	FILE			a,m,n,u,v

where

- l is any valid label (optional),
- a is the file identification number, a decimal integer in the range 1-32767,
- m is a decimal integer that defines the number of records in the file,
- n is a decimal integer in the range 1-320 that defines the length (in words) of the longest record in the file,

U is a required constant, specifying that the file must be read/written with no data conversion,

v is the associated variable, the label of a core location (variable) defined elsewhere in the program.

FILE statements must precede all other statements except HDNG, EPR, SPR, EJCT, SPAC, and LIST in the source program. The label, if used, is assigned the location of the first word of the seven words generated (see list below). The Format and Tag fields are not used and should be left blank.

Each FILE statement causes the Location Assignment Counter to be incremented by seven. The data stored in these seven words, which constitute a DEFINE FILE Table entry in the object program is as follows:

<u>Word</u>	<u>Contents</u>
1	a, the file identification number
2	m, the number of records per file
3	n, the record length (in words)
4	The address of the associated variable, v.
5	Zero. This word is filled by the Core Load Builder with the sector address of the data file. This address is relative to the address of Working Storage (with bit zero set to one) for Working Storage files and is absolute, including the drive code, for User/Fixed area files.
6	r, the number of records per sector. The number, computed by the Assembler, is the quotient of $\frac{320}{n}$ (remainder ignored)
7	b, the number of disk blocks per file. This number, computed by the Assembler, is the quotient of $\frac{16(m)}{r}$

It should be noted that the FILE statement obsoletes the \*FILE Assembler control record used with the 1130 Disk Monitor System, Version 1. Consequently, \*FILE is not recognized by the Assembler in Disk Monitor, Version 2.

## DATA DEFINITION STATEMENTS

### DMES - Define Message

The DMES statement is used to store a message within a program in a form that is acceptable to the printer output subroutines. The format of the DMES statement follows:

Label	Operation	F	T	Operands & Remarks
21	25	27	30	32 33 35 40 45 50
L	DMES		P	m

where

l is any valid label (optional),

p is the printer type code,

m is any string of valid message and control characters.

If a label is present, it is assigned to the location of the first word generated. The Tag field (column 33) is used to specify the printer type code:

<u>Tag</u>	<u>Printer</u>
b or 0	Console Printer
1	1403 Printer
2	1132 Printer

If the Tag field (printer type code) contains a character other than blank, zero, one, or two, an error results and the message is stored two EBCDIC characters per word.

The Operand field contains the control and message characters. Remarks are permitted only after an 'E' or 'b' control character.

The output generated by one DMES statement cannot exceed 60 words (120 characters). If an odd number of characters is generated, the last word is filled in with a blank, except when the statement ends with 'b'. In this case, the first character of the next DMES statement is used to fill out the word.

Control characters are used to specify certain printer operations and to define message parameters. Each control character is actually two characters, the first of which is always an apostrophe. The apostrophe (5-8 punch in IBM Card Code) is a control



delimiter and therefore is not included in the character count. The control characters and their functions or meanings are as follows:

Control Character	Function or Meaning
'X	Blank (or space)
'T	Tabulate
'D	Backspace
'B	Print black
'A	Print red
'S	Space (or blank)
'R	Carriage return
'L	Line feed
'F	Repeat following character
'E	End of message
'b	(b=blank) continues text with next DMES statement

All the above characters can be used when the printer is the Console Printer. Only 'E', 'F', 'S', 'X' and 'b' are valid control characters when the 1132 or 1403 Printer is specified; any other control characters are considered as errors.

The characters 'X' and 'S' are interchangeable. A blank character is generated for either 'X' or 'S' if the 1132 or 1403 Printer is specified; a space is generated for either 'X' or 'S' if the Console Printer is specified.

The character 'F' (repeat following character) refers only to message characters. The control characters themselves, except 'A', 'B', 'E', and 'b', can be repeated up to 99 times by inserting a number (1-99) between the apostrophe and unique control definition character. For example, '32S results in 32 space characters being inserted in the generated message.

The character 'E' is used to designate the end of the message line. The character 'b' is used to designate that the message is continued on the following DMES statement. If neither 'E' nor 'b' is included, 'E' is assumed to follow column 71. DMES statements that end with 'b' must be followed by another DMES statement.

Text apostrophes are generated by writing two successive apostrophes.

The message characters can be any valid character for the printer being used. Invalid characters are replaced with blanks.

The following example illustrates the DMES statement.

### Assembler input:

Label	Operation	F	T	Operands & Remarks
21	25	27	30	32 33 35 40 45 50 55
	DMES			'RSAMPLE PROGRAM' 'S'
	DMES			OUTPUT
	DMES			'2R'9S'1'9S2'9S3'9S4'E
	DMES			'R1234567890'123456789'
	DMES			01234567890'1234567890'E
	DMES			'2R'7X'7F'4DF(X)---
	DMES			'7X'8F'5DF'(X)---'E

### Printed output:

```

SAMPLE PROGRAM'S OUTPUT

          1          2          3          4
1234567890123456789012345678901234567890

          F(X)          F'(X)

```

Note that the device code specified in the preceding example is blank in order to generate a message for the Console Printer.

### DN — Define Name

The Define Name statement is used to convert a name specified in the Operand field of the statement to a name in Name Code in the object program. The format of this statement is shown below:

Label	Operation	F	T	Operands & Remarks
21	25	27	30	32 33 35 40 45 50
l	DN			r

where

l is any valid label (optional),

r is any valid label or name.

Name Code is truncated packed EBCDIC. The two high order bits of each character in the name are removed and the five characters are packed into the right thirty bits of two words.

```

00  C      H      A      R      S
XX|XX XXXX|XXXX XX|XX XX|XX|XXXX XX|XX XXXX

```

If a label is used, the address assigned to it is the location of the first word of the two words generated and is equal to the current value of the Location Assignment Counter. Columns 32 and 33 must be blank. The operand can have up to five characters that comply with the rules for writing symbols. The name to be converted must be left-justified in the Operand field. If remarks are used, one blank must be left between the operand and the remarks. The Location Assignment Counter is incremented by two for this statement.

## LINKING STATEMENTS

### LINK — Load Link Program

In the assembler language, the LINK statement is used to cause another core load to be loaded and executed. Only COMMON of the current core load is saved. The program loaded and executed must be specified by name. The format of the LINK statement is:

1. A symbol or blanks in the label field
2. The mnemonic, LINK, in columns 27-30
3. A valid program name in the operand field

The label of the LINK pseudo-operation is defined as the current value of the Location Assignment Counter when the LINK statement is encountered; this value is the address of the first word generated by the LINK statement.

The operand field contains a valid program name (one to five alphameric characters), left-justified in the field. The name must be present in LET/FLET at execution time. The LINK statement causes four words to be generated in the object program. The first two words contain a long BSI instruction, which branches to a specified location within the Skeleton Supervisor. The next two words contain the program name, left-justified in bits 2-32, with blanks inserted in unused rightmost positions (bits 0 and 1 are always zero). The Core Image Loader uses the core load name and begins the process required to load the new core load.

### EXIT — Return to Supervisor

In the assembler language, the EXIT statement is used to return control to the Supervisor. The format of the EXIT statement is:

1. A symbol or blanks in the label field
2. The mnemonic, EXIT, in columns 27-30

The label of the EXIT statement is defined as the current value of the Location Assignment Counter when the EXIT statement is encountered; this value is the address of the instruction generated by an EXIT statement. The operand field is ignored and can therefore be used for remarks.

The EXIT statement causes a short branch instruction to be generated in the object program. The instruction branches to a fixed location in the Skeleton Supervisor. During execution, the branch is executed and control is returned to the Supervisor. The EXIT statement should be the last logical statement in a program.

### DUMP — Dump and Terminate Execution

The DUMP statement provides an entry to the System DUMP program, which prints the contents of core storage on the principal print device in hexadecimal format.

The DUMP statement allows for flexible specification of the upper and lower limits to be dumped without altering core storage. After core has been dumped between the limits specified, the System Dump returns control to the calling program, at which point a CALL EXIT is executed. The DUMP statement is written as follows:

Label		Operation				F	T	Operands & Remarks			
21	25	27	30	32	33	35	40	45	50		
<i>l</i>		<i>DUMP</i>					<i>a<sub>1</sub></i>	<i>b<sub>1</sub></i>	<i>f</i>		

where

*l* is any valid label (optional),

*a* is any valid expression specifying the lowest-addressed core location to be dumped,

*b* is any valid expression specifying the highest-addressed core location to be dumped,

*f* is the dump format code (either blank or zero). The dump is always in hexadecimal format.

The label, if used, is assigned the location of the first of the six words generated (see list below). The Tag and Format fields must be left blank.

A DUMP statement causes the Location Assignment Counter to be incremented by six. The data stored in these six words is as follows:

Word	Contents
1 } 2 }	A long (two-word) BSI to the DUMP entry point in the Skeleton Supervisor
3	The starting address of the core dump
4	The ending location of the core dump
5	The format indicator (always zero)
6	A short branch to the EXIT entry point in the Skeleton Supervisor

If no address is specified for word 3, the dump starts in location zero. If no address is specified for word 4, the dump continues to the end of core.

A DUMP statement can be used at any point in a program; however, the user is reminded that DUMP causes a terminal DUMP to be printed. At the completion of the dump printout, the branch to EXIT is executed, thus transferring control to the Skeleton Supervisor for processing of the next job or subjob.

The format of the DUMP program output is as follows:

```
AAAA xxxx xxxx xxxx || xxxx xxxx xxxx
```

The contents (xxxx) of 16 core storage locations are printed per line. At the left is the address (AAAA) of the first location printed on that line.

**PDMP — Dump and Continue Execution**

The PDMP statement provides the ability to dump core storage between specified limits and to continue execution. The core dump is printed on the principal print device without altering core. The PDMP statement is specified in the same way as DUMP, except that PDMP appears in columns 27-30 instead of DUMP.

The PDMP statement is translated by the Assembler into a long BSI to the DUMP entry point in the Skeleton Supervisor. The parameters (operands) are converted as described in the DUMP statement (see above) except that the exit to the Supervisor is not generated for PDMP.

Upon completion of the printout of the core dump, control is returned to the next instruction following the PDMP statement to continue execution.

**LIST CONTROL STATEMENTS**

The list control statements — HDNG, LIST, SPAC, and EJCT — provide the user with the means to control and identify the assembler output listing.

**HDNG — Heading**

The HDNG statement is used to specify a one line page heading for a printed listing. The heading line consists of the data in the Operand-Remarks field.

The format of the HDNG statement is as follows.

Label	Operation	F	T	Operands & Rem	
21	25	27	30	32 33 35	40 45 50
		HDNG			PAGE HEADING

Multiple HDNG statements may be used thus allowing different sections of a listing to have different page headings.

When the 1132 or 1403 is the principal printer, the HDNG statement causes the listing to be ejected to a new page and the heading is printed. The same heading is repeated at the top of each succeeding page until a new HDNG statement is encountered.

When the Console Printer is the principal printer, the heading line is preceded by five line feeds and followed by a single line feed, and otherwise functions as a comments statement.

**LIST — List Segments of Program**

The LIST statement allows the user to list certain segments of a program on the principal printer and avoid listing other segments. The three variations of the LIST statement are shown below:

Label	Operation	F	T	Operands & Rem	
21	25	27	30	32 33 35	40 45 50
		LIST			
		LIST			ON
		LIST			OFF

The Label, Tag, and Format fields are not used with the LIST statement and should be left blank. The Operand field may be left blank or may contain the operand ON or OFF.

The LIST statement does not cause the Location Assignment Counter to be incremented.

If a LIST statement with the operand ON is encountered, the following statements, up to the next LIST statement, are listed by the Assembler.

If a LIST statement with no operand is encountered, the Assembler assumes an operand depending on the use of the LIST control record. If the LIST control record preceded the assembly, the ON operand is assumed and the Assembler acts accordingly. If the LIST control record did not precede the assembly, the OFF operand is assumed and the Assembler acts accordingly.

### SPAC — Space Listing

The SPAC statement is used to insert one or more blank lines in the listing immediately following the SPAC statement. The format of the SPAC statement is as follows:

Label	Operation	F	T	Operands & Res					
21	25	27	30	32	33	35	40	45	50
	SPAC					e			

where e is any valid positive expression.

The Label, Format, and Tag fields are not used and should be left blank.

The number of blank lines inserted in the listing is determined by the operand in the statement. The

operand can be any valid expression. The operand (expression) value must be positive; otherwise, the Assembler ignores the statement.

When the number of blank lines specified exceeds the number of lines left on the page, the page is spaced to the bottom, a restore occurs, a new heading is printed, and spacing is resumed until the number of blank lines specified has been exhausted.

The SPAC statement does not cause the Location Assignment Counter to be incremented.

### EJCT — Start New Page

The EJCT statement causes the next line of the listing to appear at the top of a new page following the page heading. The format of the EJCT statement is as follows:

Label	Operation	F	T	Operands & Res					
21	25	27	30	32	33	35	40	45	50
	EJCT								

The Label, Tag, Format, and Operand fields are not used and should be left blank.

A page overflow occurs immediately following the EJCT statement. EJCT statements may be used in succession to obtain blank pages (except for the headings printed).

The EJCT statement does not cause the Location Assignment Counter to be incremented.

## APPENDIX A. CHARACTER CODE SUMMARY

### Hexadecimal Notation

In hexadecimal notation, each digit represents a four-bit binary value. This means that a 16-bit word in the Processor-Controller can be expressed as four hexadecimal digits. The binary — hexadecimal — decimal correspondence is defined as follows:

<u>Binary</u>	<u>Hexadecimal</u>	<u>Decimal</u>
0000	0	0
0001	1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7
1000	8	8
1001	9	9
1010	A	10
1011	B	11
1100	C	12
1101	D	13
1110	E	14
1111	F	15

### Extended Binary Coded Decimal Interchange Code (EBCDIC)

In the EBCDIC code, each character is represented by a unique configuration of eight binary bits. In

the table that follows, each EBCDIC character is expressed as two hexadecimal digits.

### IBM Card Code

In the IBM Card Code, each character represents a 12-bit card-column image. In the table that follows, each card code character is expressed as four hexadecimal digits and as the card-column image.

### Paper Tape Transmission Code, 8 Channel (PTTC/8)

In the PTTC/8 code, each character is represented by a unique configuration of a case shift, plus an eight-bit code. The case shift can be common to more than one character and need be inserted only when a case shift change is necessary. In the table that follows, each character is expressed as two hexadecimal digits, followed by the case shift in parentheses.

### 1132 Printer EBCDIC Subset Hex Code

In the 1132 Printer EBCDIC subset hex code, each character is represented by a unique configuration of eight bits. In the table that follows, each 1132 Printer character is expressed as two hexadecimal digits.

### Console Printer Hex Code

In the Console Printer hexadecimal code each character is represented as two hexadecimal digits.

### 1403 Printer Hex Code

In the 1403 Printer hexadecimal code each character is represented as two hexadecimal digits.

Ref No.	EBCDIC			IBM Card Code					Graphics and Control Names	1132 Printer EBCDIC Subset Hex	PTTC/8 Hex U-Upper Case L-Lower Case	Console Printer Hex Notes	1403 Printer Hex	
	Binary		Hex	Rows										Hex
	0123	4567		12	11	0	9	8						
0	0000	0000	00	12	0	9	8	1	B030	NUL				
1	0001	0001	01	12		9		1	9010					
2	0010	0010	02	12		9		2	8810					
3	0011	0011	03	12		9		3	8410					
4	0100	0100	04	12		9		4	8210	PF	Punch Off			
5*	0101	0101	05	12		9		5	8110	HT	Horiz.Tab	41	①	
6*	0110	0110	06	12		9		6	8090	LC	Lower Case			
7*	0111	0111	07	12		9		7	8050	DEL	Delete			
8	1000	1000	08	12		9	8		8030					
9	1001	1001	09	12		9	8	1	9030					
10	1010	1010	0A	12		9	8	2	8830					
11	1011	1011	0B	12		9	8	3	8430					
12	1100	1100	0C	12		9	8	4	8230					
13	1101	1101	0D	12		9	8	5	8130					
14	1110	1110	0E	12		9	8	6	80B0					
15	1111	1111	0F	12		9	8	7	8070					
16	0001	0000	10	12	11	9	8	1	D030					
17	0001	0001	11		11	9		1	5010					
18	0010	0010	12		11	9		2	4810					
19	0011	0011	13		11	9		3	4410					
20*	0100	0100	14		11	9		4	4210	RES	Restore	05	②	
21*	0101	0101	15		11	9		5	4110	NL	New Line	81	③	
22*	0110	0110	16		11	9		6	4090	BS	Backspace	11		
23	0111	0111	17		11	9		7	4050	IDL	Idle			
24	1000	1000	18		11	9	8		4030					
25	1001	1001	19		11	9	8	1	5030					
26	1010	1010	1A		11	9	8	2	4830					
27	1011	1011	1B		11	9	8	3	4430					
28	1100	1100	1C		11	9	8	4	4230					
29	1101	1101	1D		11	9	8	5	4130					
30	1110	1110	1E		11	9	8	6	40B0					
31	1111	1111	1F		11	9	8	7	4070					
32	0010	0000	20		11	0	9	8	1	7030				
33	0001	0001	21			0	9	1	3010					
34	0010	0010	22			0	9	2	2810					
35	0011	0011	23			0	9	3	2410					
36	0100	0100	24			0	9	4	2210	BYP	Bypass			
37*	0101	0101	25			0	9	5	2110	LF	Line Feed	3D (U/L)		
38*	0110	0110	26			0	9	6	2090	EOB	End of Block	3E (U/L)	03	
39	0111	0111	27			0	9	7	2050	PRE	Prefix			
40	1000	1000	28			0	9	8	2030					
41	1001	1001	29			0	9	8	1	3030				
42	1010	1010	2A			0	9	8	2	2830				
43	1011	1011	2B			0	9	8	3	2430				
44	1100	1100	2C			0	9	8	4	2230				
45	1101	1101	2D			0	9	8	5	2130				
46	1110	1110	2E			0	9	8	6	20B0				
47	1111	1111	2F			0	9	8	7	2070				
48	0011	0000	30	12	11	0	9	8	1	F030				
49	0001	0001	31			9		1	1010					
50	0010	0010	32			9		2	0810					
51	0011	0011	33			9		3	0410					
52	0100	0100	34			9		4	0210	PN	Punch On			
53*	0101	0101	35			9		5	0110	RS	Reader Stop			
54*	0110	0110	36			9		6	0090	UC	Upper Case	09	④	
55	0111	0111	37			9		7	0050	EOT	End of Trans.			
56	1000	1000	38			9	8		0030					
57	1001	1001	39			9	8	1	1030					
58	1010	1010	3A			9	8	2	0830					
59	1011	1011	3B			9	8	3	0430					
60	1100	1100	3C			9	8	4	0230					
61	1101	1101	3D			9	8	5	0130					
62	1110	1110	3E			9	8	6	00B0					
63	1111	1111	3F			9	8	7	0070					

NOTES: Typewriter Output

- ① Tabulate
- ② Shift to black

- ③ Carrier Return
- ④ Shift to red

\* Recognized by all Conversion subroutines  
Codes that are not asterisked are recognized only by the SPEED subroutine

Ref No.	EBCDIC			IBM Card Code					Graphics and Control Names	1132 Printer EBCDIC Subset Hex	PTTC/8 Hex U-Upper Case L-Lower Case	Console Printer Hex	1403 Printer Hex		
	Binary 0123	4567	Hex	12	11	0	9	8						7-1	Hex
64*	0100	0000	40							0000	blank	40	10 (U/L)	21	7F
65		0001	41	12		0	9		1	B010					
66		0010	42	12		0	9		2	A810					
67		0011	43	12		0	9		3	A410					
68		0100	44	12		0	9		4	A210					
69		0101	45	12		0	9		5	A110					
70		0110	46	12		0	9		6	A090					
71		0111	47	12		0	9		7	A050					
72		1000	48	12		0	9		8	A030					
73		1001	49	12					8	9020					
74*		1010	4A	12					8	8820					
75*		1011	4B	12					8	8420					
76*		1100	4C	12					8	8220					
77*		1101	4D	12					8	8120					
78*	1110	4E	12					8	80A0						
79*	1111	4F	12					8	8060						
80*	0101	0000	50	12						8000	&	50	70 (L)	44	15
81		0001	51	12	11		9		1	D010					
82		0010	52	12	11		9		2	C810					
83		0011	53	12	11		9		3	C410					
84		0100	54	12	11		9		4	C210					
85		0101	55	12	11		9		5	C110					
86		0110	56	12	11		9		6	C090					
87		0111	57	12	11		9		7	C050					
88		1000	58	12	11		9		8	C030					
89		1001	59		11				8	5020					
90*		1010	5A		11				8	4820					
91*		1011	5B		11				8	4420					
92*		1100	5C		11				8	4220					
93*		1101	5D		11				8	4120					
94*	1110	5E		11				8	40A0						
95*	1111	5F		11				8	4060						
96*	0110	0000	60		11					4000	- (dash) /	60 61	40 (L) 31 (L)	84 BC	61 4C
97*		0001	61			0			1	3000					
98		0010	62		11	0	9		2	6810					
99		0011	63		11	0	9		3	6410					
100		0100	64		11	0	9		4	6210					
101		0101	65		11	0	9		5	6110					
102		0110	66		11	0	9		6	6090					
103		0111	67		11	0	9		7	6050					
104		1000	68		11	0	9		8	6030					
105		1001	69			0			8	3020					
106		1010	6A	12	11					C000					
107*		1011	6B			0			8	2420					
108*		1100	6C			0			8	2220					
109*		1101	6D			0			8	2120					
110*	1110	6E			0			8	20A0						
111*	1111	6F			0			8	2060						
112	0111	0000	70	12	11	0				E000					
113		0001	71	12	11	0	9		1	F010					
114		0010	72	12	11	0	9		2	E810					
115		0011	73	12	11	0	9		3	E410					
116		0100	74	12	11	0	9		4	E210					
117		0101	75	12	11	0	9		5	E110					
118		0110	76	12	11	0	9		6	E090					
119		0111	77	12	11	0	9		7	E050					
120		1000	78	12	11	0	9		8	E030					
121		1001	79						8	1020					
122*		1010	7A						8	0820					
123*		1011	7B						8	0420					
124*		1100	7C						8	0220					
125*		1101	7D						8	0120					
126*	1110	7E						8	00A0						
127*	1111	7F						8	0060						

Ref No.	EBCDIC			IBM Card Code					Graphics and Control Names	1132 Printer EBCDIC Subset Hex	PTTC/8 Hex U-Upper Case L-Lower Case	Console Printer Hex	1403 Printer Hex	
	Binary		Hex	Rows										Hex
	0123	4567		12	11	0	9	8						
128	1000	0000	80	12	0	8	1	B020	a b c d e f g h i					
129		0001	81	12	0		1	B000						
130		0010	82	12	0		2	A800						
131		0011	83	12	0		3	A400						
132		0100	84	12	0		4	A200						
133		0101	85	12	0		5	A100						
134		0110	86	12	0		6	A080						
135		0111	87	12	0		7	A040						
136		1000	88	12	0	8		A020						
137		1001	89	12	0	9		A010						
138		1010	8A	12	0	8	2	A820						
139		1011	8B	12	0	8	3	A420						
140		1100	8C	12	0	8	4	A220						
141		1101	8D	12	0	8	5	A120						
142		1110	8E	12	0	8	6	A0A0						
143		1111	8F	12	0	8	7	A060						
144	1001	0000	90	12	11	8	1	D020	i k l m n o p q r					
145		0001	91	12	11		1	D000						
146		0010	92	12	11		2	C800						
147		0011	93	12	11		3	C400						
148		0100	94	12	11		4	C200						
149		0101	95	12	11		5	C100						
150		0110	96	12	11		6	C080						
151		0111	97	12	11		7	C040						
152		1000	98	12	11	8		C020						
153		1001	99	12	11	9		C010						
154		1010	9A	12	11	8	2	C820						
155		1011	9B	12	11	8	3	C420						
156		1100	9C	12	11	8	4	C220						
157		1101	9D	12	11	8	5	C120						
158		1110	9E	12	11	8	6	C0A0						
159		1111	9F	12	11	8	7	C060						
160	1010	0000	A0	11	0	8	1	7020	s t u v w x y z					
161		0001	A1	11	0		1	7000						
162		0010	A2	11	0		2	6800						
163		0011	A3	11	0		3	6400						
164		0100	A4	11	0		4	6200						
165		0101	A5	11	0		5	6100						
166		0110	A6	11	0		6	6080						
167		0111	A7	11	0		7	6040						
168		1000	A8	11	0	8		6020						
169		1001	A9	11	0	9		6010						
170		1010	AA	11	0	8	2	6820						
171		1011	AB	11	0	8	3	6420						
172		1100	AC	11	0	8	4	6220						
173		1101	AD	11	0	8	5	6120						
174		1110	AE	11	0	8	6	60A0						
175		1111	AF	11	0	8	7	6060						
176	1011	0000	B0	12	11	0	8	F020						
177		0001	B1	12	11	0	1	F000						
178		0010	B2	12	11	0	2	E800						
179		0011	B3	12	11	0	3	E400						
180		0100	B4	12	11	0	4	E200						
181		0101	B5	12	11	0	5	E100						
182		0110	B6	12	11	0	6	E080						
183		0111	B7	12	11	0	7	E040						
184		1000	B8	12	11	0	8	E020						
185		1001	B9	12	11	0	9	E010						
186		1010	BA	12	11	0	8	E820						
187		1011	BB	12	11	0	8	E420						
188		1100	BC	12	11	0	8	E220						
189		1101	BD	12	11	0	8	E120						
190		1110	BE	12	11	0	8	E0A0						
191		1111	BF	12	11	0	8	E060						



Ref No.	EBCDIC		IBM Card Code				Graphics and Control Names	1132 Printer EBCDIC Subset Hex	PTTC/8 Hex U-Upper Case L-Lower Case	Console Printer Hex	1403 Printer Hex				
	Binary	Hex	12	11	0	9						8	7-1		
192	1100	0000	C0	12		0			A000	(+ zero)					
193*		0001	C1	12			1		9000	A	C1	61 (U)	3C or 3E	64	
194*		0010	C2	12			2		8800	B	C2	62 (U)	18 or 1A	25	
195*		0011	C3	12			3		8400	C	C3	73 (U)	1C or 1E	26	
196*		0100	C4	12			4		8200	D	C4	64 (U)	30 or 32	67	
197*		0101	C5	12			5		8100	E	C5	75 (U)	34 or 36	68	
198*		0110	C6	12			6		8080	F	C6	76 (U)	10 or 12	29	
199*		0111	C7	12			7		8040	G	C7	67 (U)	14 or 16	2A	
200*		1000	C8	12			8		8020	H	C8	68 (U)	24 or 26	6B	
201*		1001	C9	12			9		8010	I	C9	79 (U)	20 or 22	2C	
202		1010	CA	12	0	9	8	2	A830						
203		1011	CB	12	0	9	8	3	A430						
204		1100	CC	12	0	9	8	4	A230						
205		1101	CD	12	0	9	8	5	A130						
206		1110	CE	12	0	9	8	6	A0B0						
207		1111	CF	12	0	9	8	7	A070						
208	1101	0000	D0		11		0		6000	(- zero)					
209*		0001	D1		11			1	5000	J	D1	51 (U)	7C or 7E	58	
210*		0010	D2		11			2	4800	K	D2	52 (U)	58 or 5A	19	
211*		0011	D3		11			3	4400	L	D3	43 (U)	5C or 5E	1A	
212*		0100	D4		11			4	4200	M	D4	54 (U)	70 or 72	5B	
213*		0101	D5		11			5	4100	N	D5	45 (U)	74 or 76	1C	
214*		0110	D6		11			6	4080	O	D6	46 (U)	50 or 52	5D	
215*		0111	D7		11			7	4040	P	D7	57 (U)	54 or 56	5E	
216*		1000	D8		11			8	4020	Q	D8	58 (U)	64 or 66	1F	
217*		1001	D9		11			9	4010	R	D9	49 (U)	60 or 62	20	
218		1010	DA	12	11			9	8	2	C830				
219		1011	DB	12	11			9	8	3	C430				
220		1100	DC	12	11			9	8	4	C230				
221		1101	DD	12	11			9	8	5	C130				
222		1110	DE	12	11			9	8	6	C0B0				
223		1111	DF	12	11			9	8	7	C070				
224	1110	0000	E0			0		8	2						
225		0001	E1		11	0		9	1		2820				
226*		0010	E2			0			2		7010				
227*		0011	E3			0			3		2800	S	32 (U)	98 or 9A	
228*		0100	E4			0			4		2400	T	23 (U)	9C or 9E	
229*		0101	E5			0			5		2200	U	34 (U)	80 or B2	
230*		0110	E6			0			6		2100	V	25 (U)	B4 or B6	
231*		0111	E7			0			7		2080	W	26 (U)	90 or 92	
232*		1000	E8			0			8		2040	X	27 (U)	94 or 96	
233*		1001	E9			0			9		2020	Y	37 (U)	A4 or A6	
234		1010	EA		11	0			9	8	2010	Z	38 (U)	A0 or A2	
235		1011	EB		11	0			9	8	2		29 (U)		
236		1100	EC		11	0			9	8	3				
237		1101	ED		11	0			9	8	4				
238		1110	EE		11	0			9	8	5				
239		1111	EF		11	0			9	8	6				
240*	1111	0000	F0			0					2000	0	1A (L)	C4	49
241*		0001	F1						1		1000	1	01 (L)	FC	40
242*		0010	F2						2		0800	2	02 (L)	DB	01
243*		0011	F3						3		0400	3	13 (L)	DC	02
244*		0100	F4						4		0200	4	04 (L)	F0	43
245*		0101	F5						5		0100	5	15 (L)	F4	04
246*		0110	F6						6		0080	6	16 (L)	D0	45
247*		0111	F7						7		0040	7	07 (L)	D4	46
248*		1000	F8						8		0020	8	08 (L)	E4	07
249*		1001	F9						9		0010	9	19 (L)	E0	08
250		1010	FA	12	11	0			9	8	2				
251		1011	FB	12	11	0			9	8	3				
252		1100	FC	12	11	0			9	8	4				
253		1101	FD	12	11	0			9	8	5				
254		1110	FE	12	11	0			9	8	6				
255		1111	FF	12	11	0			9	8	7				





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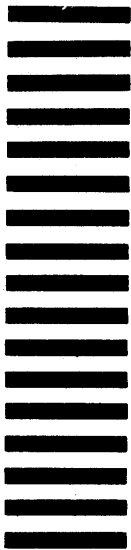
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