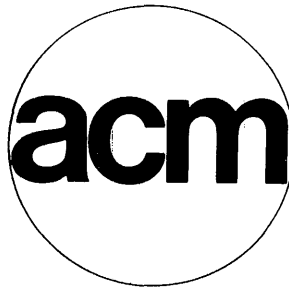


Collected Algorithms from ACM

IN THREE VOLUMES

A collation of all ACM Algorithms, including Certifications, Remarks, and Translations from the Algorithms Department of *Communications of the ACM*, 1960–1975, from *ACM Transactions on Mathematical Software*, 1975 ff, and from *ACM Transactions on Programming Languages and Systems*, 1981 ff.

Introductory Information
Volume III. Algorithms 493ff.



1981

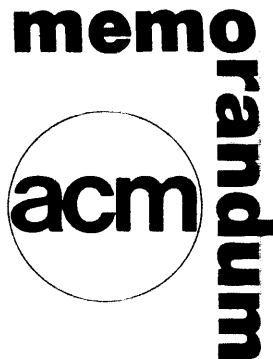
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Association for Computing Machinery

1133 AVENUE OF THE AMERICAS
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December 28, 1981

TO: Subscribers to "Collected Algorithms from ACM"

RE: Supplement No. 80

Attached is the quarterly set of sheets for your loose-leaf volume, "Collected Algorithms from ACM." The sheets in this Supplement provide the material that appeared in the algorithms section of the December 1981 issue of TRANSACTIONS ON MATHEMATICAL SOFTWARE.

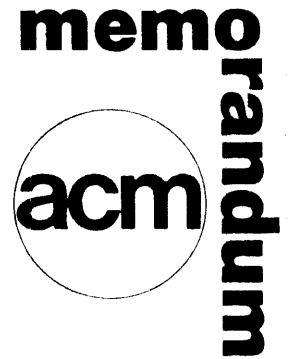
The sheets for the new algorithms 578, 579, and 580 should be added to the back of the volume in the following order:

578-P1-0 (backed by 578-P2-0)	to	578-P3-0 (backed by 578-P4-0)
579-P1-0 (backed by 579-P2-0)	to	579-P5-0 (backed by blank)
580-P1-0 (backed by 580-P2-0)		

The sheets containing the CALGO 20-year Index should be inserted with front matter in the loose-leaf binder immediately following page xiv.

Association for Computing Machinery

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April 26, 1982

TO: Subscribers to "Collected Algorithms from ACM"

RE: Supplement No. 81

Attached is the quarterly set of sheets for your loose-leaf volume, "Collected Algorithms from ACM." The sheets in this Supplement provide the material that appeared in the Algorithms Section of the March 1982 issue of TRANSACTIONS ON MATHEMATICAL SOFTWARE.

The sheets for the new algorithm 581 should be added to the back of the volume in the following order:

581-P1-0(backed by 581-P2-0) to 581-P3-0(backed by 581-P4-0)

The sheets to update algorithms previously provided should be incorporated as follows:

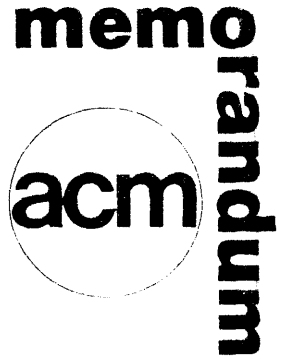
DELETE	INSERT
334-P1-R1(backed by blank)	334-P1-R1(backed by 334-P2-0)

The updated version of the "ACM Algorithms Policy" should replace the one now in the frontmatter, pages ix-xi.

The 1981 Index by Subject to Algorithms should precede the Index by Subject to Algorithms, 1960-1980(pages xv-xxvi).

Association for Computing Machinery

11 WEST 42ND STREET
NEW YORK, N.Y. 10036
(212) 869-7440



July 9, 1982

TO: Subscribers to "Collected Algorithms from ACM"

RE: Supplement No. 82

Attached is the quarterly set of sheets for your looseleaf volume, "Collected Algorithms from ACM." The sheets in this Supplement provide the material that appeared in the algorithms section of the June 1982 issue of TRANSACTIONS ON MATHEMATICAL SOFTWARE.

The sheets for the new algorithms 582, 583, and 584 should be added to the back of the volume in the following order:

582-P1-0 (backed by 582-P2-0)	to	582-P3-0 (backed by 582-P4-0)
583-P1-0 (backed by 583-P2-0)	to	583-P11-0 (backed by 583-P12-0)
584-P1-0 (backed by 584-P2-0)	to	584-P7-0 (backed by 584-P8-0)

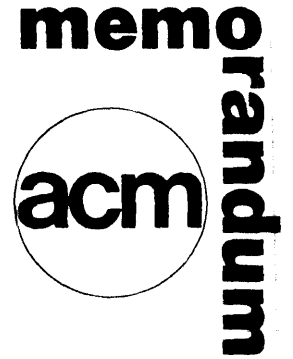
The sheets to update algorithms previously provided should be incorporated in the looseleaf binders as follows:

DELETE	INSERT
506-P11-0 (backed by blank)	506-P11-0 (backed by 506-P12-0)
508-P9-0 (backed by 508-P10-0)	508-P9-0 (backed by 508-P10-R1)
509-P9-0 (backed by 509-P10-0)	509-P9-0 (backed by 509-P10-R1)

For those with casebound volumes, the editor suggests that you insert the algorithm updates in the volume where the algorithm appears.

Association for Computing Machinery

11 WEST 42ND STREET
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October 4, 1982

TO: Subscribers to "Collected Algorithms from ACM"

RE: Supplement No. 83

Attached is the quarterly set of sheets for your looseleaf volume, "Collected Algorithms from ACM." The sheets in this Supplement provide the material that appeared in the algorithms section of the September 1982 issue of TRANSACTIONS ON MATHEMATICAL SOFTWARE.

The sheets for the new algorithms 585, 586, and 587 should be added to the back of the volume in the following order:

585-P1-0(backed by 585-P2-0)	to	585-P9-0(backed by 585-P10-0)
586-P1-0(backed by 586-P2-0)	to	586-P17-0(backed by blank)
587-P1-0(backed by 587-P2-0)	to	587-P9-0(backed by blank)

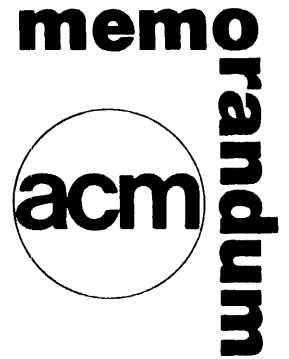
The sheets to update algorithms previously provided should be incorporated in the looseleaf binders as follows:

DELETE	INSERT
507-P9-0(backed by 507-P10-0)	507-P9-0(backed by 507-P10-R1)

For those with casebound volumes, the editor suggests that you insert the algorithm updates in the volume where the algorithm appears.

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11 WEST 42ND STREET
NEW YORK, N.Y. 10036
(212) 869-7440



January 3, 1983

To: Subscribers to "Collected Algorithms from ACM"

RE: Supplement No. 84

Attached is the quarterly set of sheets for your looseleaf volume, "Collected Algorithms from ACM." The sheets in this Supplement provide the material that appeared in the Algorithms section of the December 1982 issue of TRANSACTIONS ON MATHEMATICAL SOFTWARE.

The sheets for the algorithms 588, 589, 590, and 591 should be added to the back of the volume in the following manner:

588-P1-0(backed by 588-P2-0)	
589-P1-0(backed by 589-P2-0)	to 589-P3-0(backed by 589-P4-0)
590-P1-0(backed by 590-P2-0)	to 590-P5-0(backed by 590-P6-0)
591-P1-0(backed by 591-P2-0)	to 591-P15-0(backed by 591-P16-0)

The sheets to update algorithms previously provided should be incorporated in the looseleaf binders as follows:

DELETE	INSERT
535-P17-0(backed by 535-P18-0)	535-P17-0(backed by 535-P18-R1)
539-31-0(backed by blank)	539-P31-R1(backed by 539-P32-0)
	539-P33-0(backed by blank)
580-P1-0(580-P2-0)	580-P1-0(backed by 580-P2-R1)

The "1982 Index by Subject to Algorithms" should be inserted in front of looseleaf binder with other indexes.

For those with casebound volumes, the editor suggests that you insert the algorithm updates in the volume where the algorithm appears.

Contents

Introductory Information

Preface	vii
Algorithms Policy	ix
Software Package Policy	xiii
1980 Index	xv
1979 Index	xvi
1978 Index	xvii
1977 Index	xviii
Cumulative Index 1960–1976	xxi

Volume III. Algorithms 493ff.

Introductory Information

Preface

The Algorithms department of *Communications of the ACM* (CACM) was established in February 1960, with J. H. Wegstein as editor, for the purpose of publishing algorithms, consisting of procedures and programs, in the Algol language. In 1975 the publication of ACM algorithms material was transferred to *ACM Transactions on Mathematical Software* (TOMS) and in 1981 *ACM Transactions on Programming Languages and Systems* (TOPLAS) published its first algorithm. A wide variety of algorithms have been published and many of them have been used heavily—either in original form or as translated into other languages. Recognizing the general acceptance of the algorithm material published in CACM, TOMS, and TOPLAS the Association for Computing Machinery (ACM) has collected and reprinted the algorithms to make them more readily accessible and more serviceable to a larger group of users.

The collection contains all algorithms published in the Algorithms departments of CACM, TOMS, and TOPLAS. Covering a great variety of subjects, these algorithms include many for standard computational tasks, such as evaluation of special functions, solution of systems of linear equations, estimation of definite integrals, and sorting of data. Most of them are written in Algol or ANS Fortran.

Both to conserve space and because of changing language usage, reference material on ALGOL and Fortran, [6], [7], [10]–[13], is no longer included as it was in earlier editions. Other expositions of ALGOL are given in [2], [3], [8], and [9]. A new Fortran standard was adopted by ANSI in 1978. The official ANSI document for the new Fortran standard [14] is not yet published, but a summary is given in [4].

A cumulative index to algorithms published since 1960 is provided. In addition to algorithms published in CACM, TOMS, and TOPLAS the index lists many which appeared elsewhere. The classification scheme is a modified form of the SHARE classification. The early indexes were prepared by G. E. Forsythe (Algorithms department editor in 1964) and J. M. Varah, and succeeding editors have annually issued updated versions.

Algorithms 1–220 were originally published as received—without any refereeing whatever. Many of these have since been certified and/or corrected by their authors or by other contributors. Beginning with Algorithm 221, in March 1964, all algorithms have been refereed

independently. For many, Certifications have appeared, and modifications to some have been proposed in Remarks. In this volume, Certifications and Remarks for a given algorithm are collected with the algorithm.

Since 1964 an effort has been made to choose for publication those of the refereed algorithms that are interesting and of good quality and, at the same time, likely to be useful to others. Probably few of these algorithms would satisfy all the criteria for excellence proposed by G. E. Forsythe [5]. However, it is hoped that many will be useful and will help to disseminate good methods of solution for many problems. Some of the earlier algorithms which were not refereed are very good too, but those having Certifications and/or corrections (included in Remarks) are more likely to be valuable. It is hoped that users of those without Certifications or Remarks will contribute their experiences to TOMS or TOPLAS so that all may benefit from them. Algorithms 1–50 have been reprinted with revisions and corrections by Ageev, Alik, and Galis in the Soviet Union [1].

For general information, the present Algorithms Policy Statement of TOMS and TOPLAS is provided. The Dissemination Agreement is of particular importance: Submittal of an algorithm for publication in TOMS or TOPLAS implies that unrestricted use of the algorithm within a computer is permissible. General permission to copy the algorithm in fair use, but not for profit, is granted provided ACM's copyright notice is given and reference is made to this publication, its date of issue, and to the fact that copying is by permission of the Association for Computing Machinery.

As new Algorithms, Certifications, Remarks, and Translations appear in TOMS or TOPLAS, they are reissued quarterly to Collected Algorithms subscribers, in loose-leaf form, so that their collections may be kept up to date.

Beginning with Algorithm 493, algorithms are available in machine-readable form—tapes or cards—from ACM Algorithms Distribution Service, c/o IMSL, Sixth Floor, GNB Building, 7500 Bellaire Boulevard, Houston, TX 77063.

To facilitate the updating and to make this volume convenient to use, an understanding of the page numbering scheme for the algorithms is helpful. The page designation is in a three-part format: the left part is the

algorithm number; the middle part is the page number within the algorithm (the first page of each algorithm is P1); and the right part is the number of the revision of that page. All sheets in the original, or first, insertion of an algorithm have "0" for the right part. The first revision of a page will have a page number having the left and middle parts identical with those on the page to be replaced, but the right part will be "R1" instead of "0." The second revision of the same page would read R2, and so on. For example, 123-P2-R1 would mean the first revision of page 2 of Algorithm 123. Revised pages for an algorithm, or additional pages if required, are provided when Certifications or Remarks are added.

The Introduction that appeared in *Collected Algorithms* at its inception was written by John G. Herriot, department editor at that time. It was then revised by Fred T. Krogh, Algorithms Editor from 1976 to 1978. As present Algorithms Editor, I have prepared this revision to include current information.

Webb Miller
Algorithms Editor
Collected Algorithms from ACM

Department of Mathematics
University of California
Santa Barbara, California 93106

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11. Fortran vs. Basic Fortran. *Comm. ACM* 7 (Oct. 1964), 591-625.
12. Clarification of Fortran Standards—Initial Progress. *Comm. ACM* 12 (May 1969), 289-294.
13. Clarification of Fortran Standards—Second Report. *Comm. ACM* 14 (Oct. 1971), 628-642.
14. ANSI X3.9-1978 Programming Language FORTRAN. (Revision of X3.9-1960) to appear.

ACM Algorithms Policy

BY FRED T. KROGH¹

Matters of Content

A contribution should be in the form of an algorithm, comparison, certification, remark, or translation.

Algorithms

ACM Algorithms are published to make the fruits of software research readily available to as wide an audience as possible. An algorithm must either provide a capability not readily available or perform a task better in some way than has been done before. "Better" can mean anything from improved reliability or efficiency to more attractive packaging. In all cases, an algorithm must be of lasting value and must represent a substantial contribution in terms of the amount of work or the originality required for its creation.

In most cases the communication of new algorithmic ideas should be done either in a companion paper published in one of the ACM Transactions or in previously published work. The textual part of an algorithm submission should give a brief description of what the algorithm does and pertinent information on usage and maintenance. It should not duplicate information in another paper or in the algorithm listing. It is recognized that certain algorithms will be useful to readers of one of the ACM Transactions but deal with subjects not in the usual research domain of that journal. When appropriate, the Algorithms Editor will guide authors to other ACM journals for possible publication of such algorithms.

An ACM algorithm must be complete, portable, well documented, and well structured. The meaning of these terms is clarified below.

1. *Completeness*: With the exception of code used from a previously published ACM algorithm, a submission must include all of the code and test data necessary for the effective use and testing of the algorithm by a large segment of its intended audience. To assist those who use the algorithm, a small test driver should be provided that illustrates the use of the algorithm for a simple test problem. For testing purposes, one should provide in a single driver a sufficient variety of test cases to exercise all the main features of the code. All submitted code, including test drivers and preprocessors, is subject to the refereeing process. Code subject to more restricted use than specified in this policy may be used in a supporting role for an algorithm provided it is available from an established source for a nominal fee, and there is no nearly equivalent code available in the public domain.
2. *Portability*: It must be possible to move the code in machine readable form from one machine to another with only minor, well-documented changes. Programs written in a language having several popular dialects should use only language features common to those dialects. As initial evidence of portability, the author can either include evidence of successful execution on three different computers, i.e., computers with different basic instruction sets, or better, when available, can give the result of running the software through a verifier that checks the code for portability or conformance to a standard. All FORTRAN programs must satisfy one of the following criteria:

- (a) Conformance with the ANSI X3.9-1978 FORTRAN 77 [1].

¹ Policy revised March 1979, by Webb Miller; January 1982, by R. J. Hanson.

- (b) Pass the set of tests performed by the PFORT verifier [4].²

The Algorithms Editor will attempt to help authors who have trouble meeting either of these requirements. The PFORT verifier can be obtained at no cost from the Editor.

Machine-dependent modules, including assembly language, may be used provided:

- (a) They are clearly specified, limited in function, small, and either necessary or a substantial improvement over equivalent portable procedures.
- (b) Where possible, portable versions are also submitted to facilitate installation and testing, and to confirm the reader's understanding of the specifications. If portable versions are not provided, then tested versions for at least three different computers must be provided.
- (c) A portable test program is provided which exhaustively tests each machine-dependent module of the package.

3. *Documentation*: Each module of the code, including test drivers, must be adequately commented. Comments should include the purpose of the module, definitions of all arguments passed through the calling sequence, through global variable declarations, or obtained via input, and comments setting off and explaining major parts of the code. Comments defining machine dependencies should be gathered together and clearly flagged. Standard names and definitions for machine constants are normally to be used [2]. If these constants are used widely in the submitted algorithms, then they may be accessed through a procedure [3]. Other constants clearly identified in comments (e.g., `BIG=SQRT(SOVFLO)`) and machine dependencies not in the standard [3], may be introduced. A comment can simply point to the place where full comments are given in cases where large comment blocks would otherwise be repeated. An alphabetical list of internal variables together with how they are used is highly recommended. (Identify temporary variables as temporary.) At present, it is not possible to edit comments to provide a reference to the place in the journal where the algorithm appears. Machine-readable documentation giving more detailed information about the package than can be published is encouraged.
4. *Structure*: Code should be organized so that it is easy to use and modify, yet flexible enough to be useful for most problems in the problem area covered. Indentation should be used to identify loops, compound statements and blocks, and continuation of a statement onto another line. If integers are used for statement labels, they must be strictly increasing and should be far enough apart to allow for future modifications. If the space required by an array is highly problem dependent, then that array should have a variable dimension. (NOTE: FORTRAN requires common blocks with the same name to have the same length.)

Algorithmic ideas may be published as integral parts of regular papers without meeting these criteria. This provision may not, however, be used as a subterfuge to circumvent meeting the criteria for ACM algorithms.

Except for extremely long algorithms, all contributions will appear in their entirety in ACM's looseleaf algorithms periodical *Collected Algorithms from ACM* (CALGO). Algorithms are available in machine-readable form through the ACM Algorithms Distribution Service.

² The PFORT verifier checks a FORTRAN program for conformance with a portable subset of the former ANSI X3.9-1966 FORTRAN standard.

Comparisons

A comparison is a report on the relative merits and features of highly similar software packages for a specific subject area. This study normally includes reporting and interpreting various cogent observed facts about the packages. These facts are typically based on solving a common set of test problems. The drivers, test problems, and data used in the study are to be sent to the Algorithms Editor, at the time of submission, as an algorithm. This body of information will be used in the refereeing process and will be available as an ACM Algorithm.

Certifications

A certification is a report on a previously published algorithm. It can be a careful study of performance characteristics, a verification of correctness, or a report on extensive testing.

Remarks

A remark is a brief report on a previously published algorithm. It is usually concerned with corrections or modifications.

Translations

A translation may either provide machine-dependent modules for a machine or operating system not covered by an ACM algorithm, or a translation of the algorithm into a different high-level language. A translation will only be considered if it is a translation of an algorithm that still represents the current state of the art and it satisfies all of the criteria of a regular algorithm submission.

Submission Formalities

1. Five copies of all textual material should be sent to the Algorithms Editor. This material must be typewritten and double spaced. Companion papers to algorithms submitted to one of the ACM Transactions should be sent with the algorithm to the Algorithms Editor, as they will be refereed along with the algorithm. The purpose of those parts of the code submission that are included for testing purposes or serve only a supporting role should be clearly identified.
2. Processing of an algorithm will not proceed until it is verified that machine-readable copy of the algorithm has arrived in good form. Before sending machine-readable copy, contact the Algorithms Editor and explain the computer(s) available to you, the size of your submission, and in what form you would like to send the material. If an algorithm requires the use of code subject to restricted use, the use of this code should be cleared first with the Algorithms Editor.
3. Evidence of portability must be included with an algorithm submission.

Dissemination Agreement

Submittal of an algorithm for publication in one of the *ACM Transactions* implies that unrestricted use of the algorithm within a computer is permissible. General permission to copy and distribute the algorithm without fee is granted provided that the copies are not made or distributed for direct commercial advantage. The ACM copyright notice and the title of the publication and its date appear, and notice is given that copying is by permission of the Association for Computing Machinery. To copy otherwise, or to republish, requires a fee and/or specific permission.

Copyright Agreement

Authors of copyrightable algorithms (or their employers) are required to transfer the copyright to ACM upon acceptance of the algorithm for publication, in accordance with ACM policy to own copyright on ACM published material. The Association grants to authors the right to reuse their material and seeks the permission of authors before approving republication.

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Software Package Policy

The ACM Algorithm Distribution Service makes it possible to remove the length limitation on algorithms published in TOMS and TOPLAS and allows them to publish and disseminate large programs.

Technically, a collection of algorithms is just another algorithm, but practically ACM editors will distinguish between “ordinary” algorithms and larger entities called polyalgorithms, software packages, systemized collections, etc. This policy statement is an extension of the algorithms policy of TOMS and TOPLAS which describes those software packages that these Transactions are interested in publishing and gives the specific requirements expected for such packages.

General Criteria and Procedures

Usefulness is an important criterion for the acceptance of any ACM algorithm; it is an essential criterion for ACM software packages. An acceptable ACM package must address itself to a standard and widely occurring problem set. Furthermore, it must be designed to ease the user’s task as much as possible consistent with the nature of the problem set. The human engineering aspects of the package will be weighted heavily in its evaluation by the editors.

All ACM algorithms, but especially packages, are expected to be of high quality in all respects. ACM packages must use good algorithms in their components, they must use good programming (including style, structure, and adherence to standards), they must be well documented (both for use and modification), and they must be thoroughly tested and evaluated.

Transportability will also be weighted heavily in that it directly affects usefulness. A wonderful package will not be accepted if it requires a specific environment of machine, operating system and I/O devices that is available only at a very few locations. However, complete transportability will not be insisted upon, and in each case a judgment will be made by the editors on whether

a reasonable and useful level of transportability has been achieved. It is widely realized that, in general, high level language algorithms cannot yet be constructed to give top performance in a variety of environments.

ACM publication normally consists of two parts: a description of the package and a critical review of it. Authors of a package will submit material for the first part and the editors will seek referee/reviewers for the second part. We note that a multitude of widely distributed software packages already exist, some of excellent quality. We consider such packages to have already been published. Other packages may become widely distributed in the future. Authors independent of the developers of these packages are invited to submit critical reviews of them. An author of a new package is expected to demonstrate the superiority of the new package over any existing packages.

Specific Requirements

The following guidelines for authors give information about the expected contents of papers to be printed in TOMS or TOPLAS, the nature of the packages to be distributed, and the material to be submitted for consideration by the editors.

(1) *The Descriptive Paper.* The paper should contain a clear statement of goals, intended users, expected environment, and the problem domain of the package. Then there should be a discussion of the main components and steps in the development of the package. This includes mathematical descriptions, general algorithm structure, testing and evaluation, documentation, programming, and portability. Interesting design questions should be discussed here. It is not expected that a complete mathematical derivation be given; rather one should give pertinent references. Descriptions of new techniques belong in an independent paper. The aim here is to give the reader a general and reasonably complete description without burdening him with the full details of documentation and techniques. Full details are available from the references and the ACM Algorithm Distribution Service. The paper should give some basic statistics about the package, e.g. length, examples of performance (memory usage, execution times), I/O requirements.

This policy was originally written for TOMS in 1975 by Editor-in-Chief John R. Rice, with Richard M. Heiberger and Charles L. Lawson, who made substantial contributions to the formulation of the policy, and with Edward Battiste, David C. Hoaglin, and Paul Velleman, who made valuable suggestions. This policy now applies also to TOPLAS, in which the first algorithm was published in 1981.

(2) *The Package.* The ACM Algorithm Distribution Service makes three items available: the programs, user documentation, and internal documentation. All documentation is incorporated into the program as machine-readable comments. A package must be of acceptable quality in all of the following aspects: user documentation; internal documentation; ease of use; programming style; consistency of style, structure, and notation; basic algorithmic and mathematical techniques; adherence to language standards; ease of modification and transportation; performance characteristics; overall structure and organization.

(3) *Material To Be Submitted.* The material for an ACM software package should be submitted to the Algorithms Editor or to the Editor-in-Chief of TOMS or of TOPLAS and should consist of the following items:

- machine readable form of package
- descriptive paper (three copies)
- documentation (user and internal—probably part of first item)
- drivers for examples and testing
- test and evaluation results (listings from actual runs, summaries, etc.).

The drivers submitted should exercise all the modules in the package. The test and evaluation results must support the claim of superiority over existing packages; a separate discussion of this aspect will probably be required.

It must be emphasized that finding referees for software packages is likely to be difficult, and that if competent referees cannot be found it will be necessary to reject the manuscript. Authors will improve their chances of finding willing referees if they submit their material in a neat and well organized form. All computer tests submitted should be carefully annotated so that the referee can easily determine the nature of the test, the test data used, the expected results, and the results produced by the package.

(4) *The Critical Review.* This paper will analyze the strengths and weaknesses of the package in an objective manner. It may include extensive discussions of only a few aspects of the package if they are of sufficient interest (particularly for already existing packages). Critical reviews of widely disseminated packages not published in either TOMS or TOPLAS must contain basic information about their distribution, listings, and maintenance support.

Editorial Comment

It is clear that the preparation of a publishable software package is a task well beyond the usual technical publication. The critical review part of publication is a way in which referees can obtain suitable recognition for the substantial effort involved in a thorough evaluation of a package.

1982 Index by Subject to Algorithms

Algorithms published elsewhere are included here. A cumulative index to algorithms published in *Communications*, the *ACM Transactions on Mathematical Software*, and the *ACM Transactions on Programming Languages and Systems* from 1960 through 1980 may be consulted in "Collected Algorithms from ACM," which is a looseleaf collection of all algorithms, each with certifications and remarks, if any, appended. The 1981 index appears in the December 1981 issue of *Communications*.

QUADRATURE [D1]

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ALL OTHERS [Z]

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Index By Subject To Algorithms, 1960–1980

Algorithms are grouped according to the classification system below, which is a modification of the SHARE classification system. Algorithms published elsewhere are included; see list of abbreviations for other sources.

Each entry gives the key to the classification system categories (A1, B1, . . .); the number of the algorithm if it exists; abbreviated title with language designator (A for Algol, F for Fortran, etc.—these appear on the more recent algorithms); source designator, volume, year, page (in parenthesis).

Classification System			
Z	All Others	B4	Roots and Powers
S	Approximation of Special Functions	MS	Scaling and Conversion of Data
G6	Combinations and Permutations	C6	Series, Summation and Convergence Acceleration of
L2	Compiling	G1	Simple Calculations on Statistical Data
A2	Complex Arithmetic	O2	Simulation of Computing Structure
I5	Composite Input	F4	Simultaneous Linear Equations
O2	Computing Structure Simulation	E3	Smoothing
C6	Convergence Acceleration	M1	Sorting
M2	Conversion and Scaling of Data	S	Special Functions, Approximation of
G2	Correlation and Regression Analysis	G2	Statistical Data, Correlation and Regression Analysis of
E2	Curve and Surface Fitting	G1	Statistical Data, Simple Calculations on
M2	Data Conversion and Scaling	G7	Subset Generators
F3	Determinants	C6	Summation of Series
D2	Differential Equations, Ordinary	E2	Surface and Curve Fitting
D3	Differential Equations, Partial	R2	Symbol Manipulation
D4	Differentiation	B1	Trig and Inverse Trig Functions
F2	Eigenvalues and Eigenvectors of a Matrix	F5	Vectors, Orthogonalization of
B3	Exponential and Logarithmic Functions	C5	Zeroes of one or more Nonlinear Equations
E4	Function Minimizing and Maximizing	C2	Zeroes of Polynomials
H	Graph Structures, Operations Research		
B2	Hyperbolic Functions		Classification of Special Functions
I5	Input—Composite	S04	Bernoulli and Euler Numbers and Polynomials
D5	Integral Equations	S18	Bessel Function, Modified
E1	Interpolation	S19	Bessel Functions of Complex Argument
F1	Inversion of a Matrix	S18	Bessel Functions of Pure Imaginary Argument
F4	Linear Equations, Simultaneous	S17	Bessel Functions of Real Argument
B3	Logarithmic and Exponential Functions	S20	Bessel and Related Functions, Miscellaneous
F2	Matrix Eigenvalues and Eigenvectors	S14	Beta Function and Incomplete Beta Function
F1	Matrix Operations, Including Inversion	S03	Binomial Coefficients
F3	Matrix, Determinant of	S07	Circular Functions, Miscellaneous
E4	Maximizing a Function	S19	Complex Argument, Bessel Functions of
E4	Minimizing a Function	S13	Cosine Integrals
C5	Nonlinear Equations, Zeroes of	S23	Curve-Fitting
A1	Number Theory	S04	Derivatives and Differences of Zero
C1	Operations on Polynomials and Power Series	S15	Derivatives
H	Operations Research, Graph Structures	S04	Differences and Derivatives of Zero
D2	Ordinary Differential Equations	S23	Differentiation, Numerical
F5	Orthogonalization	S21	Elliptic Integrals and Functions
D3	Partial Differential Equations	S15	Error Integral
G6	Permutations and Combinations	S04	Euler and Bernoulli Numbers and Polynomials
Y1	Physics Applications	S13	Exponential Integrals
J6	Plotting	S14	Factorial Function
C1	Polynomials, Operations on	S03	Factorials
C2	Polynomials, Zeroes of	S22	Functions: Miscellaneous Higher Mathematical Functions
C1	Power Series, Operations on	S14	Gamma Function and Incomplete Gamma Function
B4	Powers and Roots	S15	Hermite Polynomials and Functions
D1	Quadrature	S15	Higher Integrals
G5	Random Number Generators	S22	Higher Mathematical Functions, Miscellaneous
A1	Real Arithmetic	S18	Imaginary Argument, Bessel Functions of
G2	Regression and Correlation	S14	Incomplete Beta and Gamma Functions
K2	Relocation	S13	Integrals of Exponentials, Logarithms, Sines, Cosines, etc.
		S21	Integrals, Elliptic

S15	Integrals: Higher Integrals and the Error Integral	S03	Partitions
S23	Integration, Numerical	S14	Polygamma Function
S23	Interpolation	S15	Polynomials, Hermite
S04	Inverse Powers, Sums of	S04	Powers and Inverse Powers, Sums of
S19	Kelvin Functions	S14	Psi Function
S16	Legendre Functions	S13	Sine Integrals
S13	Logarithmic Integrals	S07	Spherical Functions, Miscellaneous
S22	Miscellaneous Higher Mathematical Functions	S04	Sums of Powers and of Inverse Powers
S18	Modified Bessel Functions	S21	Theta Functions
S15	Moments	S07	Trigonometric Functions, Natural
S23	Numerical Differentiation and Integration	S04	Zero, Differences and Derivatives of

CLASSIFICATION SYSTEM (MODIFIED SHARE)

A1	REAL ARITHMETIC, NUMBER THEORY
A2	COMPLEX ARITHMETIC
B1	TRIG AND INVERSE TRIG FUNCTIONS
B2	HYPERBOLIC FUNCTIONS
B3	EXPONENTIAL AND LOGARITHMIC FUNCTIONS
B4	ROOTS AND POWERS
C1	OPERATIONS ON POLYNOMIALS AND POWER SERIES
C2	ZEROS OF POLYNOMIALS
C5	ZEROS OF ONE OR MORE NON-LINEAR EQUATIONS
C6	SUMMATION OF SERIES, CONVERGENCE ACCELERATION
D1	QUADRATURE
D2	ORDINARY DIFFERENTIAL EQUATIONS
D3	PARTIAL DIFFERENTIAL EQUATIONS
D4	DIFFERENTIATION
D5	INTEGRAL EQUATIONS
E1	INTERPOLATION
E2	CURVE AND SURFACE FITTING
E3	SMOOTHING
E4	MINIMIZING OR MAXIMIZING A FUNCTION
F1	MATRIX OPERATIONS, INCLUDING INVERSION
F2	EIGENVALUES AND EIGENVECTORS OF MATRICES
F3	DETERMINANTS
F4	SIMULTANEOUS LINEAR EQUATIONS
F5	ORTHOGONALIZATION
G1	SIMPLE CALCULATIONS ON STATISTICAL DATA
G2	CORRELATION AND REGRESSION ANALYSIS
G5	RANDOM NUMBER GENERATORS
G6	PERMUTATIONS AND COMBINATIONS
G7	SUBSET GENERATORS
H	OPERATIONS RESEARCH, GRAPH STRUCTURES
I5	INPUT - COMPOSITE
J6	PLOTTING
K2	RELOCATION
L2	COMPILING
M1	SORTING
M2	DATA CONVERSION AND SCALING
O2	SIMULATION OF COMPUTING STRUCTURE
R2	SYMBOL MANIPULATION
S	APPROXIMATION OF SPECIAL FUNCTIONS...
S	FUNCTIONS ARE CLASSIFIED S01 TO S22, FOLLOWING
S	FLETCHER-MILLER-ROSENHEAD, INDEX OF MATH. TABLES
Y1	PHYSICS APPLICATIONS
Z	ALL OTHERS

ABBREVIATIONS FOR SOURCES

AL	ARGONNE NATIONAL LABORATORY, NATS PROJECT
AS	APPLIED STATISTICS
BT	B. I. T.
CA	COMMUNICATIONS OF THE ASSOCIATION FOR COMPUTING MACHINERY
CB	COMPUTER BULLETIN
CC	COMPUTER PHYSICS COMMUNICATIONS
CH	CHIFFRES
CJ	COMPUTER JOURNAL
CP	COMPUTING
HA	U. K. A. E. A., HARWELL, REPORTS
IC	I. C. C. BULLETIN (INTERNATIONAL COMPUTATION CENTRE)
JC	JOURNAL OF THE ASSOCIATION FOR COMPUTING MACHINERY
MC	MATHEMATICS OF COMPUTATION
NM	NUMERISCHE MATHEMATIK
SN	SIAM JOURNAL ON NUMERICAL ANALYSIS
TO	TRANSACTIONS ON MATHEMATICAL SOFTWARE
ZM	ZASTOSOWANIA MATEMATYKI

LANGUAGE CODE

A	ALGOL
B	BASIC
F	FORTRAN
M	ASSEMBLY LANGUAGE
P	PL/1
PF	PL/1 - FORMAC

ACM
INDEX

- A1 REAL ARITHMETIC, NUMBER THEORY
- A1 PARTITION FCNS. (MODULO D) BT, 9-69(83)
- A1 RATIONAL INTERVAL FUNCTIONS (A) BT, 14-74(87)
- A1 7 EUCLIDEAN ALGORITHM CA, 3-60(240)
- A1 35 SIEVE OF ERATOSTHENES CA, 4-61(151), 5-62(209), 5-62(438), 10-67(570)
- A1 35 CA, 4-61(319)
- A1 61 RANGE ARITHMETIC CA, 4-61(339), 4-61(498)
- A1 68 AUGMENTATION CA, 4-61(498), 5-62(439)
- A1 72 COMPOSITIONS CA, 5-62(344), 5-62(514)
- A1 93 GENERALIZED ARITHMETIC CA, 5-62(344)
- A1 95 PARTITIONS CA, 5-62(345), 5-62(557)
- A1 99 JACOBI SYMBOL CA, 5-62(434)
- A1 114 PARTITIONS CA, 5-62(556), 8-65(170)
- A1 139 DIOPHANTINE EQUATION CA, 7-64(243)
- A1 223 PRIME TWINS CA, 7-64(481), 7-64(702)
- A1 237 GREATEST COMMON DIVISOR CA, 8-65(493)
- A1 262 RESTRICTED PARTITIONS OF N CA, 8-65(493)
- A1 263 PARTITION GENERATOR CA, 8-65(493)
- A1 264 MAP OF PARTITIONS INTO INTEGERS CA, 10-67(451), 11-68(14)
- A1 307 SYMMETRIC GROUP CHARACTERS CA, 10-67(569), 10-67(570)
- A1 310 PRIME NUMBER GENERATOR 1 13-70(192)
- A1 310 CA, 10-67(570), 10-67(570)
- A1 311 PRIME NUMBER GENERATOR 2 CA, 10-67(666)
- A1 313 MULTI-DIMENSION PARTITION GEN. CA, 12-69(563)
- A1 356 PRIME NUMBER GENERATOR CA, 12-69(563)
- A1 357 PRIME NUMBER GENERATOR CA, 13-70(52)
- A1 371 PARTITIONS IN NATURAL ORDER CA, 13-70(52), 13-70(695)
- A1 372 COMPLEX PRIMES CA, 13-70(120)
- A1 374 RESTRICTED PARTITIONS GEN. CA, 13-70(120)
- A1 375 DOUBLY RESTRICTED PARTITIONS CA, 13-70(447), 16-73(257)
- A1 386 GCD OF N INTEGERS & MULTPLRS(F) CA, 13-70(693), 13-70(695)
- A1 401 COMPLEX PRIMES CA, 14-71(48)
- A1 403 CIRCULAR INTEGER PARTITION CA, 16-73(379)
- A1 448 MULTIPLY-RESTRICTED PARTITNS(F) CA, 16-73(699)
- A1 469 ARITH. OVER FINITE FIELD (A) C.J., 9-67(416)
- A1 18 SUM OF FACTORS OF N (SUMFAC) C.J., 11-68(347), 12-69(293)
- A1 25 RATIONAL APPROX. OF REAL NO. C.J., 13-70(282)
- A1 PARTITIONING INTEGERS, N DIM C.J., 14-71(166)
- A1 FACTORING ALGORITHM C.J., 14-71(447), 15-72(285)
- A1 71 17-74(379)
- A1 72 MULTIPLE INTEGER ARITHMETIC (A) C.J., 15-72(281)
- A1 SIMULATION OF INTERVAL ARITH. CP, 5-70(377)
- A1 APPROX OF SUMS OF MANY TERMS(A) CP, 12-74(323)
- A1 EXTENSION OF ARITH. PRECISION NM, 18-71(224)
- A1 524 MULTIPLE-PRECISION ARITHMETIC (F) TO, 4-78(71)
- A1 524 MULTIPLE-PRECISION ARITHMETIC (F) TO, 4-78(71), 5-79()
- A1 8 TO SHORTEN SEQ OF INTEGER SETS ZM, 11-70(357)
- A2 COMPLEX ARITHMETIC
- A2 COMPLEX ARITHMETIC BT, 2-62(233)
- A2 116 COMPLEX DIVIDE CA, 5-62(435)
- A2 186 COMPLEX ARITHMETIC CA, 6-63(386)
- A2 312 COMPLEX ABS. SORT CA, 10-67(665)
- A2 19 ADD, SUB, MULT, DIVD-COMPLEX C.J., 10-67(112), 10-67(208)
- A2 86 COMPLEX INTERVAL ARITHMETIC (F) C.J., 18-75(83)
- B1 TRIG AND INVERSE TRIG FUNCTIONS
- B1 ARCCOS(Z) BT, 2-62(236)
- B1 ARCSIN(Z) BT, 2-62(236)
- B1 ARCTAN(Z) BT, 2-62(236)
- B1 206 ARCCOSSIN CA, 6-63(519), 8-65(104)
- B1 229 ELEMENTARY FCNS. BY CONT. FRACT. CA, 7-64(296), 12-69(692)
- B1 241 ARCTAN(Z) CA, 7-64(546)
- B1 FOURIER TRANSFORM CC, 2-71(127)
- B1 COS FCN. BY CHEBYSHEV EXPANSION NM, 4-62(411), 7-65(195)
- B1 SIN FCN. BY CHEBYSHEV EXPANSION NM, 4-62(411), 7-65(194)
- B1 ARCSIN BY CHEBYSHEV EXPANSION NM, 4-62(412)
- B1 ARCTAN BY CHEBYSHEV EXPANSION NM, 4-62(412)
- B1 TAN FCN. BY CHEBYSHEV EXPANSION NM, 4-62(412), 7-65(195)
- B1 7 EVALUATION OF A TRIG POLY ZM, 11-70(353)
- B2 HYPERBOLIC FUNCTIONS
- B2 COSH(X) BT, 2-62(235)
- B2 SINH(X) BT, 2-62(235)
- B3 EXPONENTIAL AND LOGARITHMIC FUNCTIONS
- B3 NIELSENS GENERALIZED POLYLOG BT, 10-70(38)
- B3 46 EXP(Z), Z COMPLEX CA, 4-61(178), 5-62(347)
- B3 48 LOG(Z), Z COMPLEX CA, 4-61(179), 5-62(347), 5-62(391), 7-64(485)
- B3 48 CA, 7-64(660), 8-65(279)
- B3 243 LOGARITHM OF COMPLEX NUMBER CA, 7-64(660), 8-65(279)
- B3 EXP FCN. BY CHEBYSHEV EXPANSION NM, 4-62(410)
- B3 LOG FCN. BY CHEBYSHEV EXPANSION NM, 4-62(411)
- B4 ROOTS AND POWERS
- B4 53 ROOTS OF COMPLEX NUMBERS CA, 4-61(180), 4-61(322)
- B4 106 POWERS OF COMPLEX NUMBER CA, 5-62(388), 5-62(557)
- B4 190 POWERS OF COMPLEX NUMBERS CA, 6-63(388)
- C1 OPERATIONS ON POLYNOMIALS AND POWER SERIES
- C1 29 POLYNOMIAL SHIFTER CA, 3-60(604)
- C1 131 DIVIDE POWER SERIES CA, 5-62(551)
- C1 134 EXPONENTIATE POWER SERIES CA, 5-62(553), 6-63(390)
- C1 158 EXPONENTIATE POWER SERIES CA, 6-63(104), 6-63(390), 6-63(522)
- C1 158
- C1 193 REVERT POWER SERIES CA, 6-63(388), 6-63(745)
- C1 273 SOLN. OF EQNS. BY REVERSION CA, 9-66(11)
- C1 305 SYMMETRIC POLYNOMIALS CA, 10-67(450), 11-68(272)
- C1 337 POLY. AND DERIV. BY HORNER SCHEME CA, 11-68(633), 12-69(39)
- C1 446 CHEBYSHEV SERIES (F) CA, 16-73(254), 18-75(276)
- C1 EVALUATION OF CONTINUED FRACTNS CH, 9-XX(327)
- C1 15 CALCULATION OF GRAM POLYS C.J., 9-66(323)
- C1 498 AIRY FNS USING CHEBYSHEV SERIES APPR. (F) TO, 1-75(372)
- C2 ZEROS OF POLYNOMIALS
- C2 11 LEHMERS METHOD ET, 4-64(255)
- C2 21 ROTATING-CROSS METHOD BT, 7-67(244)
- C2 NEWTON METHOD (A) BT, 13-73(71)
- C2 3 BAIRSTOW CA, 3-60(74), 3-60(354), 4-61(105), 4-61(153), 4-61(181)
- C2 3
- C2 30 BAIRSTOW-NEWTON CA, 3-60(643), 4-61(238), 5-62(50), 10-67(293)
- C2 30
- C2 59 RESULTANT METHOD CA, 4-61(236)
- C2 75 RATIONAL ROOTS-INTEG. COEFF. CA, 5-62(48), 5-62(392), 5-62(439)
- C2 78 RATIONAL ROOTS-INTEG. COEFF. CA, 5-62(97), 5-62(168), 5-62(440)
- C2 78
- C2 105 NEWTON-MAEHLI CA, 5-62(387), 6-63(389)
- C2 174 BOUNDS ON ZEROS CA, 6-63(311)
- C2 256 MODIFIED GRAEFFE METHOD CA, 8-65(379), 9-66(687)
- C2 283 REAL SIMPLE ROOTS CA, 9-66(273)
- C2 326 ROOTS OF LOW ORDER POLY EQNS CA, 11-68(269)
- C2 340 RT-SQUARING AND RESULTANT METH. CA, 11-68(779), 12-69(281)
- C2 419 ZEROS OF COMPLEX POLYNOMIAL (F) CA, 15-72(97), 17-74(157)
- C2 429 ROOTS OF A POLYNOMIAL (F) CA, 15-72(776), 16-73(579)
- C2 21 BAIRSTOW C.J., 10-67(207)
- C2 SOLUTION OF POLY. TO MACH. PRECISION (A) C.J., 18-75(258)
- C2 ZEROS OF POLYNOMIALS CP, 5-70(109)
- C2 APP VALUES OF POLY ROOTS CP, 6-70()
- C2 PL/1 PROG FOR LEHMERS METHOD MC, 23-69(829)
- C2 493 ZEROS OF A REAL POLYNOMIAL (F) TO, 1-75(178)
- C2 4 POLY DECOMP INTO QUAD FACTORS ZM, 11-70(215)
- C2 6 HOMOGRAPHIC TRANSF OF ZEROS ZM, 11-70(229)
- C2 1963(364)
- C3 553 EXPLICIT TIME INTEG., PARABOLIC EQNS (F) TO, 6-80(236)
- C5 ZEROS OF ONE OR MORE NON-LINEAR EQUATIONS
- C5 6 ZEROS BY INTERP. OR BISECTION BT, 3-63(205)
- C5 MODIFIED REGULA FALSI METHOD BT, 11-71(168)
- C5 ZEROS OF ENTIRE FUNCTIONS (A) BT, 13-73(8)
- C5 HIGH ORDER REGULA FALSI (A) BT, 13-73(253)
- C5 SOLVE F(X)=0 BT, 17-77(179)
- C5 ROOT ISOLATION USING FUNCTION VALS (F) BT, 18-78(311)
- C5 2 REGULA FALSI CA, 3-60(74), 3-60(354), 3-60(475), 4-61(153)
- C5 2
- C5 4 BISECTION CA, 3-60(174), 4-61(153)
- C5 15 REGULA FALSI CA, 3-60(475), 3-60(602), 4-61(153)
- C5 15
- C5 25 REAL ZEROS CA, 3-60(602), 4-61(153), 4-61(154)
- C5 25
- C5 26 REGULA FALSI CA, 3-60(603), 4-61(153)
- C5 107 SOLN. OF SIMULTANEOUS NONLINEAR EQNS (F) CP, 22-79(282)
- C5 110 NONLINEAR EQUATIONS WITH A PARAMETER (F) CP, 23-80(85)
- C5 196 MULLERS METHOD CA, 6-63(442), 11-68(12)
- C5 314 N FUNCTIONAL EQNS. IN N UNKNOWN CA, 10-67(726), 12-69(38)
- C5 315 DAMPED TAYLR SERIES-NONLIN. SYS. CA, 10-67(726), 12-69(513)
- C5 316 NON-LINEAR SYSTEM CA, 10-67(728), 14-71(493)
- C5 365 COMPLEX ROOT FINDING CA, 12-69(686)
- C5 378 SOLVE NONLINEAR SYSTEM OF EQS. CA, 13-70(259)
- C5 413 TAYLOR COEFF. OF ANALYTIC FUN CA, 14-71(669)
- C5 443 SOLUTION OF W*EXP(W)=X (F) CA, 16-73(123), 17-74(225)
- C5 554 SYSTEMS OF NONLINEAR EQUATIONS (F) TO, 6-80(240)
- C5 555 FIXED POINTS OR ZEROS OF C2 MAPS (F) TO, 6-80(252)
- C5 44 SOLN. OF NONLINEAR EQUATIONS C.J., 12-69(406), 13-70(219)
- C5 BRACKETING TECHNIQUE C.J., 13-70(101)
- C5 ALG WITH GUARANTEED CONVERGENCE C.J., 14-71(422)
- C5 ROOT F(T)=0 WITH ERROR BOUNDS CP, 2-67(231)
- C5 ROOTS OF FCNS. WITH ERROR BOUND CP, 4-69(197)
- C5 12 SOLN. FIN. DIM. NONLINEAR SYSTEM CP, 5-70(84)
- C5 ENCLOSING ZEROS OF A FUNCTION CP, 5-70(356)
- C5 SYSTEMS OF NONLINEAR EQUATIONS CP, 8-71(41)
- C5 REAL ROOTS IN AN INTERVAL (A) CP, 9-72(327)
- C5 SYSTEM OF NON-LINEAR ALG EQ (F) HA, AERE-R5947
- C5 SPARSE SYSTEM NON-LINEAR EQ (F) HA, AERE-R7293
- C5 BROUWER DEGREE IN R*R (A) MC, 27-73(133)
- C5 SYSTEMATIC SEARCH IN HIGH DIM. SETS (F) SN, 14-77(296)
- C5 502 DEP. SOL. OF NONLIN. SYS. ON A PARAM. (F) TO, 2-76(98)
- C6 SUMMATION OF SERIES, CONVERGENCE ACCELERATION
- C6 83 COMPLEX DISCRETE FAST FOURIER TRANSFORM (F) AS, 24-75(153)
- C6 97 REAL FAST FOURIER TRANSFORM (F) AS, 25-76(166)
- C6 117 CHIRP DISCRETE FOURIER TRANSF. (F) AS, 26-77(351)

- C6 1 FIND LIMIT OF SEQUENCE BT,1-61(64)
C6 EPSILON ALGORITHM BT,2-62(240)
C6 LOW ROUND-OFF SUMMATION METHOD BT,11-71(271)
C6 8 EULER SUM CA,3-60(311),6-63(663)
C6 128 FOURIER SERIES SUMMATION CA,5-62(513),7-64(421)
C6 157 FOURIER SERIES APPROXIMATION CA,6-63(103),6-63(521),
C6 157 6-63(618)
C6 215 EPSILON ALGORITHM CA,6-63(662),7-64(297)
C6 255 FOURIER COEFFICIENTS CA,8-65(279),12-69(636)
C6 277 CHEBYSHEV SERIES COEFFICIENTS CA,9-66(86)
C6 320 HARMONIC ANAL-SYM DISTR DATA CA,11-68(114)
C6 338 FAST FOURIER TRANSFORM CA,11-68(773)
C6 339 FAST FOURIER WITH ARB. FACTORS CA,11-68(776),12-69(187)
C6 345 CONVOLUTION BASED ON FFT CA,12-69(179),12-69(566)
C6 393 SUMMATION WITH ARB. PRECISION CA,13-70(570),15-72(468)
C6 473 LEGENDRE SERIES COEFFS (F) CA,17-74(25)
C6 EPSILON ALG.-CONTINUED FRACTNS CH,9-XX(327)
C6 SUM FOURIER SERIES CJ,6-63(248)
C6 31 COMPLEX FOURIER ANALYSIS CJ,10-68(414),11-68(115)
C6 GENZD. EULER TRANSFORMATION CJ,14-71(437)
C6 UNSCRAMBLE FAST FOURIER TRANS. CP,3-68(334)
C6 38 CONVERGENCE ACCELERATION (F) CP,21-78(87)
C6 EPSILON ALGORITHM NM,6-64(22)
C6 545 OPTIMIZED MASS STORAGE FFT (F) TO,5-79()
- D1 QUADRATURE
D1 4 AUX FN FOR DISTR. INTEGRALS (F) AS,17-68(190),22-73(428)
D1 2 ADAPTIVE SIMPSONS RULE BT,1-61(290)
D1 8 ROMBERG METHOD BT,4-64(58)
D1 1 QUADRATURE CA,3-60(74)
D1 32 MULTIPLE INTEGRAL CA,4-61(106),6-63(69),
D1 32 11-68(826)
D1 60 ROMBERG METHOD CA,4-61(255),5-62(168),
D1 60 5-62(281),7-64(420)
D1 84 SIMPSONS RULE CA,5-62(208),5-62(392),
D1 84 5-62(440),5-62(557)
D1 98 COMPLEX LINE INTEGRAL CA,5-62(345)
D1 103 SIMPSONS RULE CA,5-62(347)
D1 125 GAUSSIAN COEFFICIENTS CA,5-62(510)
D1 145 ADAPTIVE SIMPSON CA,5-62(604),6-63(167),
D1 145 8-63(171)
D1 146 MULTIPLE INTEGRAL CA,5-62(604),7-64(296)
D1 182 ADAPTIVE SIMPSON CA,6-63(315),7-64(244)
D1 198 ADAPTIVE MULTIPLE INTEGRAL CA,6-63(443)
D1 233 MULTIPLE INTEG.-SIMPSONS RULE CA,7-64(348),13-70(512)
D1 257 HAVIE: INTEGRATOR CA,8-65(381),9-66(795),
D1 257 9-66(871)
D1 279 CHEBYSHEV QUADRATURE CA,9-66(270),9-66(434),
D1 279 10-67(294),10-67(666)
D1 280 GREGORY QUADRATURE COEFFICIENTS CA,9-66(271)
D1 281 ROMBERG QUADRATURE COEFFICIENTS CA,9-66(271),10-67(188)
D1 303 ADAPTIVE QUAD.-RANDM PANEL SIZE CA,10-67(373)
D1 331 GAUSSIAN QUADRATURE FORMULAS CA,11-68(432),12-69(280)
D1 331 13-70(512)
D1 351 MODIFIED ROMBERG QUADRATURE CA,12-69(324),13-70(374)
D1 351 13-70(263),13-70(449)
D1 353 FILON QUADRATURE CA,12-69(457),13-70(263)
D1 379 ADAPTIVE SIMPSON QUADRATURE CA,13-70(260),15-72(107)
D1 400 MODIFIED HAVIE INTEGRATION (F) CA,13-70(622),17-74(324)
D1 417 WEIGHTS OF INTERPOLATION QUAD CA,14-71(807)
D1 418 FOURIER INTEGRAL (F) CA,15-72(47),15-72(469)
D1 418 16-73(775),17-74(324)
D1 424 CLENSHAW-CURTIS QUADRATURE CA,15-72(353)
D1 424 CLENSHAW-CURTIS QUADRATURE (F) CA,15-72(353),5-79(240)
D1 427 FOURIER COSINE INTEGRAL CA,15-72(358)
D1 436 PRODUCT TYPE TRAPEZOIDAL RULE CA,15-72(1070)
D1 437 PRODUCT TYPE SIMPSONS RULE CA,15-72(1070)
D1 438 GAUSS-LEGENDRE-SIMPSON RULE CA,15-72(1071)
D1 439 GAUSS-LEGENDRE-SIMPSON RULE CA,15-72(1072)
D1 440 MULTI-DIM MONTE CARLO QUAD. (A) CA,16-73(49)
D1 453 GAUSS QUAD FORM BRWCH'S INT(F) CA,16-73(486)
D1 468 INTEGR. OVER FINITE INTERVAL(F) CA,16-73(694)
D1 MONTE CARLO QUADRATURE CJ,6-63(281)
D1 48 MSR ANALYSIS INTEGRAL CJ,13-70(207),14-71(215)
D1 CLENSHAW-CURTIS QUADRATURE CJ,15-72(141)
D1 CHEBYSHEV SERIES FOR OSCILL. INTEG. (A) CJ,19-76(258)
D1 101 CAUTIOUS ADAPTIVE QUADRATURE (A) CJ,21-78(180)
D1 QUADRATURE WITH ERROR BOUNDS CP,3-68(47)
D1 FOUR-DIM. ROMBERG INTEGRATION CP,9-72(45)
D1 22 INT OF HIGHLY OSCILLATORY FN(F) CP,13-74(183)
D1 32 AUTO. COMT. OF SINGULAR INTEG. (F) CP,17-76(265)
D1 33 WLF-QUADRATURE (F) CP,81-77(271)
D1 34 INTEGRATION OVER A TRIANGLE (F) CP,19-77(179)
D1 36 NUMERICAL INTEGRATION (F) CP,20-78(363)
D1 CALC. OF GAUSS QUAD. RULES MC,23-69(221)
D1 OPTIMAL ADDIT'N OF ABSISSAS(F) MC,28-74(344)
D1 ROMBERG METHOD NM,6-64(15)
D1 QUADRATURE BY EXTRAPOLATION NM,9-66(274)
D1 CENTROID METHOD INTEGRATION NM,16-71(343)
D1 GAUSSIAN QUADRATURE NM,18-72(465)
D1 43 EVAL. IMPROPER INTEGRAL ON FINITE INT. (A)ZM,15-76(141)
- D2 9 RUNGE-KUTTA CA,3-60(312),9-66(273)
D2 194 ZEROS OF O.D.E. SYSTEM CA,6-63(441)
D2 218 KUTTA-MERSON CA,6-63(737),7-64(585),
D2 218 9-66(273)
D2 407 DIFSUB FOR SOLUTION OF ODE CA,14-71(185)
- D2 461 CUBIC SPLINE APPROXIMATION (F) JA,16-73(635)
D2 SOLN. BY TAYLOR SERIES METHOD CJ,14-71(243)
D2 77 AUTOMATIC STEP CHANGE METHOD(A) CJ,16-73(187)
D2 INTEGR. SYSTEMS WITH DISCONTINUITIES (F) CJ,17-74(275)
D2 SOL. OF LIN. SECOND-ORDER DE(A) CP,13-74(143)
D2 EXTRAPOLATION METHOD NM,8-66(10)
D2 1ST ORDER, AUTO. STEP CHANGE NM,14-70(317)
D2 497 AUTO INTEG. OF FUNCTIONAL DIFF EQNS (F) TO,1-75(369)
D2 504 GERK: GLOBAL ERROR EST. FOR ODES (F) TO,2-76(200)
D2 534 STIFF DIFFERENTIAL EQUATIONS (F) TO,4-78()
D2 47 SOL. 1-ST ORD. ODE'S IN CHEBY. SERIES(A) ZM,15-76(252)
- D3 PARTIAL DIFFERENTIAL EQUATIONS
D3 CONFORMAL MAP-ELLIPSE TO CIRCLE HT,2-62(243)
D3 392 SYSTEMS OF HYPERBOLIC PDE CA,13-70(567),15-72(107)
D3 460 OPTIM. PARAM. ADI PROCEDURES(F) CA,16-73(633)
D3 DELSQPHI, TWO-DIM POISSON EQ CC,2-71(139)
D3 POISSON EQ IN CYLIN COORD CC,2-71(157)
D3 FINITE DIFFERENCE METHODS (M) CC,4-72(82)
D3 SOL OF LAPLACE'S EQUATION (F) HA,T.P.422
D3 A FAST CAUCHY-RIEMANN SOLVER (F) MC,33-79(585)
D3 KERNEL FCN. IN BNDY.VALUE PROBS. NM,3-61(209)
D3 BNDY.VALUE PROBS.-INTEGRAL OPRS NM,7-65(56)
D3 225(37)NO.24
D3 METHOD OF CHARACTERISTICS FOR PDES (A) NM,24-75(331)
D3 HYPERBOLIC DIFFERENTIAL EQUATIONS (A) NM,28-77(121)
D3 HYPERBOLIC DIFF EQNS OF SECOND ORDER (A) NM,34-80(217)
D3 494 SOLUTIONS OF SYSTEMS OF PDES (F) TO,1-75(261)
D3 527 GENERALIZED MARCHING ALGORITHM (F) TO,4-78(165)
D3 540 GENERAL COLLOCATION PDE SOFTWARE (F) TO,5-79(326)
D3 541 SEPERABLE ELLIPTIC PDE'S (F) TO,5-79(352),5-79(365)
D3 543 FAST SOLNS. OF HELMHOLTZ-TYPE PDE'S (F) TO,5-79()
D3 48 BOUND. VAL. PROB. FOR BIHARMONIC EQ. (A) ZM,15-76(397)
- D4 DIFFERENTIATION
D4 INDICES IN FAA DI BRUNO FMLA(F) BT,13-73(38)
D4 79 DIFFERENCE EXPRESSION COEFF. CA,5-62(97),6-63(104)
D4 DIFFN. BY NEVILLES FORMULAS NM,8-66(462)
- D5 INTEGRAL EQUATIONS
D5 368 INVERSION OF LAPLACE TRANSFORM CA,13-70(47),13-70(624)
D5 486 NUM. INVERSION OF LAPLACE TRANSFORM CA,17-74(587),2-76(395)
D5 486 NUM. INVERSION OF LAPLACE TRANSFORM (A) CA,17-74(587),2-76(395)
D5 486 3-77(111)
D5 INVERSION OF ABEL'S INTEGRAL EQ. (F) CC,10-75(98)
D5 CALCULATION OF GREEN'S FUNCTIONS (F) CC,11-76(27)
D5 INTEGR. OF UNEQ. SPACED DATA (A) CJ,15-72(81)
D5 97 SOLN. OF LINEAR FREDHOLM EQUATIONS (A) CJ,20-77(374)
D5 INVERSION LAPLACE TRANSFORM. CP,2-67(153)
D5 1965(933)
D5 FREDHOLM INTEGRAL EQUATIONS (A) NM,34-80(125)
D5 503 FREDHOLM EQS. OF SECOND KIND (F) TO,2-76(196)
- E1 INTERPOLATION
E1 51 LOG-LINEAR FIT (F) AS,21-72(218),25-76(193)
E1 52 POWER SUMS OF DEVIATIONS AS,21-72(226)
E1 SMOOTH CURVE INTERPOLATION ET,9-69(69)
E1 18 RATIONAL INTERP.-CONT.FRACT. CA,3-60(508),5-62(437)
E1 70 AITKEN INTERPOLATION CA,4-61(497),5-62(392)
E1 77 INTERPOLATION, DIFFN., INTEGRN. CA,5-62(96),5-62(348),
E1 77 6-63(446),6-63(663)
E1 167 CONFLUENT DIVIDED DIFFERENCES CA,6-63(164),6-63(523)
E1 168 INTERPOLATION-DIVIDED DIFFCES. CA,6-63(165),6-63(523)
E1 169 INTERPOLATION-DIVIDED DIFFCES. CA,6-63(165),6-63(523)
E1 187 DIFFCES.AND DERIVS.-RECURSIVE CA,6-63(387)
E1 210 LAGRANGE INTERPOLATION CA,6-63(616)
E1 211 HERMITE INTERPOLATION CA,6-63(617),6-63(619)
E1 264 INTERPOLATION IN A TABLE CA,8-65(602)
E1 416 COEFF OF INTERPOLATION FORMULA CA,14-71(806)
E1 472 INTERPOLATING NATURAL SPLINE(A) CA,16-73(763)
E1 480 SMOOTH & INT. NATURAL SPLINE(A) CA,17-74(463)
E1 10 AITKEN INTERPOLATION CJ,9-66(211)
E1 NEVILLE INTERPOLATION CJ,9-66(212)
E1 40 SPLINE INTERPOLN OF DEGREE 3 CJ,12-69(198),12-69(409)
E1 42 QUINTIC SPLINES INTERPOLATION CJ,12-69(292),13-70(115)
E1 62 INTERPOLATE QUINTIC SPLINES CJ,13-70(437)
E1 CUBIC INTERPOLATION CJ,15-72(80)
E1 73 PERIODIC CUBIC SPLINE (A) CJ,15-72(282)
E1 74 PERIODIC SPLINE INTERPOLAT'N(A) CJ,15-72(283),17-74(91)
E1 DISCRETE APPROXIMATION (A) CJ,16-73(180)
E1 90 PERIODIC CUBIC SPLINE EQUIDIST. NODES(A) CJ,18-75(183)
E1 10 TWO-DIM. SMOOTH INTERPOLATION CP,4-69(180),8-71(200)
E1 EXP. SPLINE INTERPOLATION CP,4-69(230)
E1 HIGH DEGREE LIDSTONE SPLINES CP,7-71(65)
E1 QUINTIC SPLINES CP,7-71(75)
E1 16 TWO-DIM. EXPONENTIAL SPLINES CP,7-71(365)
E1 TRIGONOMETRIC INTERPOLATION (A) CP,16-76(319)
E1 CUBIC SPLINE APPROX. TO FN (F) HA,AERE-R7308
E1 COMPUTATION OF EXPONENTIAL SPLINE (A) NM,35-80(81)
E1 NONLINEAR SPLINE FUNCTIONS (F) SN,14-77(254)
E1 507 QUINTIC NATURAL SPLINE INTERPOLATION (F) TO,2-76(281)
E1 526 BIVARIATE INTERP. AND SURFACE FITTING (F)TO,4-78(161),5-79(242)
E1 526 BIVARIATE INTERP. AND SURFACE FITTING (F)TO,4-78(161)
E1 547 CUBIC SPLINE INTERP. AND SMOOTHING (F) TO,6-80(92)
E1 2 CALCULATION OF RATIONAL FUNCT ZM,11-70(101)
E1 10 BEST POLYNOMIAL FOR REMEZ ALG. ZM,12-71(107)
E1 46 INTERP. QUADRATIC SPLINE FUNCTION (A) ZM,15-76(245)
- D2 ORDINARY DIFFERENTIAL EQUATIONS
D2 194 ZEROS OF O.D.E. SYSTEM CA,3-60(312),9-66(273)
D2 218 KUTTA-MERSON CA,6-63(441)
D2 218 CA,6-63(737),7-64(585),
D2 218 9-66(273)
D2 407 DIFSUB FOR SOLUTION OF ODE CA,14-71(185)

- E2 CURVE AND SURFACE FITTING
- E2 50 TESTS OF FIT AS, 21-72 (103)
- E2 74 LI-NORM FIT OF A STRAIGHT LINE (F) AS, 23-74 (244), 25-76 (96)
- E2 110 L SUB P NORM FIT OF A STRAIGHT LINE (F) AS, 26-77 (114), 28-79 (112)
- E2 110 L SUB P NORM FIT OF A STRAIGHT LINE (F) AS, 26-77 (114)
- E2 CONTINUED FRACTION EXPANSION BT, 2-62 (245)
- E2 28 LEAST SQUARES BY ORTHOG. POLYN CA, 3-60 (604), 4-61 (544), 10-67 (293)
- E2 28
- E2 37 ECONOMICIZATION CA, 4-61 (151), 5-62 (438), 6-63 (445)
- E2 37
- E2 38 ECONOMICIZATION CA, 4-61 (151), 6-63 (445)
- E2 74 LEAST SQUARES WITH CONSTRAINTS CA, 5-62 (47), 6-63 (316)
- E2 91 CHEBYSHEV FIT CA, 5-62 (281), 6-63 (167), 7-64 (290), 10-67 (803)
- E2 91
- E2 164 SURFACE FIT CA, 6-63 (162), 6-63 (450)
- E2 176 SURFACE FIT CA, 6-63 (313), 15-72 (1073)
- E2 177 LEAST SQUARES WITH CONSTRAINTS CA, 6-63 (313), 6-63 (390)
- E2 275 EXPONENTIAL CURVE FIT CA, 9-66 (85)
- E2 276 CONSTRAINED EXPONENTIAL FIT CA, 9-66 (85)
- E2 295 EXPONENTIAL CURVE FIT CA, 10-67 (87)
- E2 296 LEAST SQ. FIT-ORTHOG. POLYS. CA, 10-67 (87), 10-67 (377)
- E2 296
- E2 318 CHEBYSHEV CURVE FIT (REVISED) 12-69 (636)
- E2 375 FIT DATA TO A*EXP(-B*X) CA, 10-67 (801)
- E2 376 FIT DATA TO A*COS(B*X+C) CA, 13-70 (120)
- E2 409 DISCRETE CHEBYSHEV CURVE FIT CA, 13-70 (120)
- E2 414 CHEBYSHEV APP OF CONT FUNCTION CA, 14-71 (739)
- E2 433 SMOOTH CURVE FITTING (F) CA, 15-72 (914), 2-76 (208)
- E2 458 INTERVAL LINEAR PROGRAMMING (F) CA, 16-73 (629)
- E2 474 BIVARIATE INTERPOLATION (F) CA, 17-74 (26)
- E2 474 BIVARIATE INTERPOLATION (F) CA, 17-74 (26), 5-79 (241)
- E2 476 CURVE FITTING USING SPLINES (F) CA, 17-74 (220)
- E2 485 G-SPLINES VIA FACTORIZATION (F) CA, 17-74 (526)
- E2 PARAMETER SEARCH CC, 1-69 (135)
- E2 MULTIVARIATE LEAST-SQUARES (F) CC, 7-74 (56)
- E2 OFF-DIAGONAL RATIONAL APPROXIMANTS (F, A) CC, 9-75 (46)
- E2 33 EXPONENTIAL FIT CJ, 11-68 (114)
- E2 CURVE-FITTING PROGRAM (F) CC, 11-76 (211)
- E2 37 EXPONENTIALLY DAMPED LINEAR FIT CJ, 12-69 (100)
- E2 67 AXIS INVARIANT PROCEDURE CJ, 14-71 (209), 14-71 (451)
- E2 69 TRIGONOMETRIC CURVE FITTING CJ, 14-71 (213), 15-72 (79)
- E2 85 GENERATION OF CHISHOLM APPROXIMANTS (A) CJ, 18-75 (81)
- E2 AN ALGORITHM FOR DRAWING F(X,Y)=0 (A) CJ, 19-76 (246)
- E2 42 CUBIC SPLINES WITH CONVEXITY CONSTR. (F) CP, 24-80 (349)
- E2 CURVE THRU SEQUENCE OF PTS (F) HA, AERE-R7092
- E2 RATIONAL CHEBYSHEV APPROX. JC, 11-64 (66)
- E2 LI APPROX. ON A DISCRETE SET NM, 8-66 (299)
- E2 CHEBYSHEV APPROX.-DISCRETE SET NM, 8-66 (303)
- E2 RATIONAL CHEBYSHEV APPROX. NM, 9-66 (177)
- E2 REMES ALGORITHM-GENERALIZED NM, 10-67 (203)
- E2 RATIONAL CHEBYSHEV APPROX. NM, 10-67 (291)
- E2 RATIONAL CHEBYSHEV APPROX. NM, 12-68 (242)
- E2 RATIONAL FUNCTION APPROXIMATION NM, 18-71 (127)
- E2 PACKAGE FOR CALCULATING WITH B-SPLINES (F) SN, 14-77 (441)
- E2 501 TRANS. ALG. 409, DISCRETE CHEBY. FIT (F) TO, 2-76 (95)
- E2 501 DISCRETE CHEBYSHEV CURVE FIT (F) TO, 2-76 (95), 4-78 (95)
- E2 510 PIECEWISE LINEAR APPROX. TO TAB. DATA (F) TO, 2-76 (388)
- E2 514 CUBIC CURVE FITTING USING LOCAL DATA (A) TO, 3-77 (175)
- E2 525 ADAPTIVE SMOOTH CURVE FITTING (F) TO, 4-78 (82)
- E2 32 CONSTRUCTION OF POLY. VALUES (A) ZM, 14-74 (327), 14-75 (646)
- E3 SMOOTHING
- E3 73 CROSS-SPECTRUM SMOOTHING (F) AS, 23-74 (238)
- E3 SMOOTHING WITH SPLINE FUNC. BT, 10-70 (501)
- E3 188 SMOOTHING CA, 6-63 (387)
- E3 189 SMOOTHING CA, 6-63 (387)
- E3 216 SMOOTHING CA, 6-63 (663)
- E3 SMOOTHING BY SPLINE FCNS. NM, 10-67 (203)
- E3 547 CUBIC SPLINE INTERP. AND SMOOTHING (F) TO, 6-80 (92)
- E4 MINIMIZING OR MAXIMIZING A FUNCTION
- E4 47 SIMPLEX METHOD FOR FUNCTION MINIMUM (F) AS, 20-71 (338), 23-74 (250)
- E4 47 23-74 (252)
- E4 47 SIMPLEX METHOD FOR FUNCTION MINIMUM (F) AS, 20-71 (338), 23-74 (250)
- E4 47 25-76 (97), 27-78 (380)
- E4 47 SIMPLEX METHOD FOR FUNCTION MIN. (F) AS, 20-71 (388), 25-76 (97)
- E4 133 OPTIMUM OF ONE-DIM. MULTIMODAL FUN. (F) AS, 27-78 (366)
- E4 STRICT ESTIMATION OF MAXIMUM BT, 11-71 (199)
- E4 129 MINIMIZE FUNCT. OF N VARIABLES CA, 5-62 (550), 6-63 (521)
- E4 178 DIRECT SEARCH CA, 6-63 (313), 9-66 (684), 11-68 (498), 12-69 (637), 12-69 (638)
- E4 178
- E4 203 MINIMIZE FUNCT. OF N VARIABLES CA, 6-63 (517), 7-64 (585), 8-65 (171)
- E4 203
- E4 204 MINIMIZE FUNCT. OF N VARIABLES CA, 6-63 (519)
- E4 205 MINIMIZE FUNCT. OF N VARIABLES CA, 6-63 (519), 8-65 (171)
- E4 251 FUNCTION MINIMIZATION CA, 8-65 (169), 9-66 (686), 12-69 (512), 14-71 (358)
- E4 251
- E4 315 MINIMIZING SUM OF SQUARES CA, 10-67 (726), 12-69 (513)
- E4 387 FCN. MINIMIZATION CA, 13-70 (509)
- E4 450 ROSENBERG FUNCTION MINIMIZATION (F) CA, 16-73 (482), 17-74 (590)
- E4 450
- E4 454 CMLX METH CONSTRAINED OPTIM (F) CA, 16-73 (487), 17-74 (471)
- E4 2 FIBONACCI SEARCH CB, 8-64 (147), 9-65 (105), CB, 9-65 (104), CJ, 9-67 (41)
- E4 7 MIN. OF UNIMODAL FCN. OF 1 VAR. CJ, 10-75 (343)
- E4 FN. MIN. & ANAL. OF PARAM. ERRORS (F) CJ, 7-64 (151)
- E4 MINIMIZING FCN.-CONJ. GRAD. CJ, 9-67 (414), 9-67 (416)
- E4 2
- E4 16 MINIMIZING ARG. OF UNIMODAL FCN. CJ, 9-67 (415)
- E4 17 MIN. OF A UNIMODAL FCN OF 1 VAR. CJ, 9-67 (415)
- E4 46 MODIFIED DAVIDON METHOD CJ, 13-70 (111), 14-71 (106)
- E4 46 14-71 (214)
- E4 87 MINIMUM OF NON-LINEAR FUNCTION (F) CJ, 18-75 (86)
- E4 109 GOLDEN SECTION SEARCH (A) CJ, 22-79 (383)
- E4 MIN AND MAX OF POLYNOMIAL CP, 6-70 (35)
- E4 GEN. QUADRATIC PROGRAMMING (F) HA, AERE-R6370
- E4 UNCONSTRAINED MINIMIZATION (F) HA, AERE-R6469
- E4 MIN BY CONJUGATE GRADIENTS (F) HA, AERE-R7073
- E4 MIN BY QUASI-NEWTON METHODS (F) HA, AERE-R7125
- E4 500 MIN. OF UNCONSTRAINED MULTIVAR. FUN. (F) TO, 2-76 (87)
- E4 500 MIN. OF UNCONSTRAINED MULTIVAR. FUN. (F) TO, 2-76 (87), 2-76 (397)
- E4 500 3-77 (112)
- E4 500 MIN. OF UNCONSTRAINED MULTIVAR. FUN. (F) TO, 2-76 (87), 6-80 ()
- E4 559 STATIONARY PT OF QUAD FUN, LIN CONST (F) TO, 6-80 (432)
- E4 12 MIN OF FUNCTION OF ONE VARIABLE ZM, 12-71 (221)
- E4 30 MIN OF FUNCTION OF ONE VAR (A) ZM, 14-74 (137)
- E4 31 MIN OF FUN. OF SEVERAL VARS (A) ZM, 14-74 (143)
- F1 MATRIX OPERATIONS, INCLUDING INVERSION
- F1 6 TRIANGULAR DECOMP OF SYMM MATRIX (F) AS, 17-68 (195), 18-69 (118)
- F1 6 23-74 (477)
- F1 6 TRIANGULAR DECOMP OF SYMM MATRIX (F) AS, 17-68 (195), 18-69 (118)
- F1 6 23-74 (477), 27-78 (309)
- F1 7 POS SEMI-DEF SYM MAT INVERSION AS, 17-68 (198), 18-69 (118)
- F1 11 SYMMETRIC MATRIX NORMALIZATION AS, 17-68 (287)
- F1 12 SUM OF SQUARE AND PRODUCT MAT AS, 17-68 (289)
- F1 34 SEQ INVERSION OF BAND MATRIX AS, 19-70 (290)
- F1 37 INVERSION OF SYMM. MATRIX (A) AS, 20-71 (111), 23-74 (100)
- F1 60 LATEST ROOTS & VECT SYMM MAT (F) AS, 22-73 (260), 23-74 (101)
- F1 127 GEN. OF RANDOM ORTHOGONAL MATRICES (F) AS, 27-78 (199)
- F1 20 SMITH NORMAL FORM BT, 7-67 (163)
- F1 26 SINGULAR VALUE DECOMPOSITION BT, 10-70 (376)
- F1 42 INVERSION CA, 4-61 (176), 4-61 (498), 6-63 (38), 6-63 (445)
- F1 50 INVERSE OF HILBERT MATRIX CA, 4-61 (179), 5-62 (50), 6-63 (38)
- F1 50
- F1 51 INVERSE OF PERTURBED MATRIX CA, 4-61 (180), 5-62 (391)
- F1 52 INVERSE OF TEST MATRIX CA, 4-61 (180), 4-61 (339), 5-62 (438), 6-63 (39), 6-63 (446)
- F1 58 INVERSION-GAUSSIAN ELIMINATION CA, 4-61 (236), 5-62 (347), 5-62 (438), 5-62 (606)
- F1 66 INVERSION-SQRT METHOD CA, 4-61 (322), 5-62 (50), 5-62 (348)
- F1 67 GRAM MATRIX CA, 4-61 (322), 5-62 (348)
- F1 120 INVERSION-GAUSSIAN ELIMINATION CA, 5-62 (437), 6-63 (40), 6-63 (445)
- F1 140 INVERSION CA, 5-62 (556), 6-63 (448)
- F1 150 INVERSE OF SYMMETRIC MATRIX CA, 6-63 (67), 6-63 (390), 7-64 (148)
- F1 150
- F1 166 MONTE CARLO INVERSE CA, 6-63 (164), 6-63 (523)
- F1 197 MATRIX DIVISION CA, 6-63 (443), 5-62 (606)
- F1 230 MATRIX PERMUTATION CA, 7-64 (347)
- F1 231 INVERSION-GAUSS. ELIM.-COMP. PIV. CA, 7-64 (347), 8-65 (220)
- F1 274 HILBERT DERIVED TEST MATRIX CA, 9-66 (11), 12-69 (407)
- F1 287 INVERSE MATRIX TRIANGULATION CA, 9-66 (513)
- F1 298 SQ. RT. OF A POS. DEFINITE MATRIX CA, 10-67 (182), 12-69 (325)
- F1 319 TRIANG FCTRS OF MODIFIED MATRIX CA, 11-68 (12)
- F1 325 ADJUST INVERSE OF SYM MATRIX CA, 11-68 (118)
- F1 348 MTRX SCALING BY INTGR PRGRMNG CA, 12-69 (212)
- F1 358 SING. VAL. DECOMP.-COMPLEX MATRIX CA, 12-69 (564)
- F1 380 TRANSPOSE OF RECTANGULAR MATRIX CA, 13-70 (325), 13-70 (327)
- F1 380 15-72 (49)
- F1 467 TRANSPOSITION IN PLACE (F) CA, 16-73 (692)
- F1 467 TRANSPOSITION IN PLACE (F) CA, 16-73 (692), 5-79 ()
- F1 MATRIX PADE APPROXIMANTS (F) CP, 11-76 (325)
- F1 12 INVERSION-SYMM POS DEF MAT CJ, 9-66 (321)
- F1 20 PERMUTATIONS OF ROWS AND COLS. CJ, 10-67 (206)
- F1 53 DECOMP OF POS DEF SYM MTRX CJ, 13-70 (421)
- F1 70 VAR. BANDWIDTH POS. DEF. EQU. CJ, 14-71 (446)
- F1 SYMM. DECOMP. OF POS. DEF. BAND MTRX CP, 1-66 (77)
- F1 ADJUST INVERSE OF SYM. MATRIX CP, 3-68 (76)
- F1 14 SYMMETRIC MATRIX CP, 7-71 (127)
- F1 15 INVERSION OF TRIANGULAR MATRIX CP, 7-71 (327)
- F1 METHOD OF COMPLETING CP, 8-71 (221)
- F1 INVERT SYMM BUT NOT POS DEF MATRIX (A) CP, 14-75 (131)
- F1 INVERT RECTANG MAT FULL RANK (F) HA, AERE-R6072
- F1 INVERSION AND DETERMINANT (F) HA, AERE-R6899
- F1 EQUIVALENCE OF MATRICES IC, 3-64 (62)
- F1 SYMMETRIC MATRIX IN MAX-NORM JC, 18-71 (566)
- F1 MANIPULATION OF TENSOR PROD (A) MC, 27-73 (595)
- F1 UPDATE OF GRAM-SCHMIDT QR FACT. (A) MC, 30-76 (772)
- F1 INVERSE-CONFL. VANDERMONDE MTRX NM, 5-63 (429)
- F1 SYMM. DECOMP. OF POS. DEF. BAND MTRX NM, 7-65 (357)
- F1 SYMM. DECOMP. OF POS. DEF. MTRX NM, 7-65 (368)
- F1 HOUSEHOLDER TRIDIAG OF SYM MTRX NM, 11-68 (184)
- F1 TRIDIAG. OF SYM. BAND MATRIX NM, 12-68 (234)
- F1 SIM. RED. GEN. MTRX. HESSENBERG FORM NM, 12-68 (354)
- F1 SINGULAR VALUE DECOMPOSITION NM, 14-70 (403)
- F1 BAND SYMMETRIC MATRIX NM, 16-70 (85)
- F1 MAT MULT AND TRIANGULAR DECOMP NM, 16-70 (145)
- F1 SYMMETRIC MATRICES NM, 16-70 (205)
- F1 REDUCING SUBSPACES BY BLOCK DIAGON. (F) SN, 16-79 (359)
- F1 508 MATRIX BANDWIDTH & PROFILE REDUCTION (F) TO, 2-76 (375)
- F1 509 HYBRID PROFILE REDUCTION (F) TO, 2-76 (378)
- F1 513 ANALYSIS OF IN-SITU TRANSPOSITION (F) TO, 3-77 (104), 5-79 ()
- F1 513 ANALYSIS OF IN-SITU TRANSPOSITION (F) TO, 3-77 (104)
- F1 529 PERMUTATIONS TO BLOCK TRIANG. FORM (F) TO, 4-78 (189)
- F1 539 BASIC LINEAR ALGEBRA SUBPROGRAMS (F) TO, 5-79 (324)

- F2 EIGENVALUES AND EIGENVECTORS OF MATRICES
 F2 EIGENSYSTEM PACKAGE, EISPACK(F) AL,2-72,6-72
 F2 SYMMETRIC-BISECTION, INV. ITN. BT,4-64(124)
 F2 COMPUTING SU(N) INVARIANTS BT,11-71(1)
 F2 SYMMETRIC TRIDIAGONAL MATRIX BT,11-71(262)
 F2 85 JACOBI METHOD CA,5-62(208),5-62(440),
 F2 85 6-63(447)
 F2 104 REDUCTION-BAND TO TRIDIAGONAL CA,5-62(387)
 F2 122 GIVENS TRIDIAGONAL REDUCTION CA,5-62(482),7-64(144)
 F2 183 REDUCTION-BAND TO TRIDIAGONAL CA,6-63(315)
 F2 253 SYMMETRIC QR-EIGENVALUES CA,8-65(217),10-67(376)
 F2 254 SYMMETRIC QR-EIGENVALUES, EIVECTORS CA,8-65(218),10-67(376)
 F2 270 EIGENVECTORS BY GAUSSIAN ELM. CA,8-65(668)
 F2 297 SYM. SYS. (A-LAM=B)X, EIVALS-VECS. CA,10-67(181)
 F2 343 REAL GENERAL MATRIX CA,11-68(820),13-70(122)
 F2 343 13-70(694),15-72(466)
 F2 384 EIVAL-EIVEC OF A SYMM. MATRIX CA,13-70(369),13-70(750)
 F2 405 ROOTS OF MATRIX PENCILS CA,14-71(113),15-72(107)
 F2 464 REAL, SYMM, TRIDIAGONAL MAT. (A) CA,16-73(689)
 F2 HERMITIAN MAT EIGENPROBLEM $H^*X=H^*X$ (F) CC,8-74(85)
 F2 BANDED EIGENVALUE PROBLEMS (F) CC,10-75(30)
 F2 TRIDIAGONAL SIMIL. BY ELM. CJ,4-61(175)
 F2 EIGENVALUES BY QR-ALGORITHM CJ,4-62(344)
 F2 SYMM MAT-LIT AND STURM SEQ. CJ,9-66(303)
 F2 32 EIVALS-REAL SYM MAT-DBL QR STP CJ,11-68(112)
 F2 LR ALGRTHM FOR SYMM TRID MAT(A) CJ,15-72(268)
 F2 JK METHOD FOR REAL SYMM MAT (F) CJ,15-72(271)
 F2 EIGENVALUES OF QUINDIAGONAL MATRIX (A) CJ,18-75(70)
 F2 93 GENERALIZED EIGENVALUE PROBLEM (F) CJ,20-77(86)
 F2 EIGENVALUES-LAGUERRES METHOD MC,18-64(474),20-66(437)
 F2 EIVECS FOR GENL COMPLEX MAT MC,22-68(785)
 F2 EIGENVALUES OF TRIDIAG. MATRIX NM,4-62(354)
 F2 EIGENVECTORS OF TRIDIAG. MTRX NM,4-62(354)
 F2 HOUSEHOLDERS METHOD NM,4-62(354)
 F2 LR TRANSFORMATION METHOD NM,5-63(273)
 F2 HOUSEHOLDER RED.-COMPLEX MAT. NM,8-66(79)
 F2 JACOBI METHOD NM,9-66()
 F2 EIGENVECTORS OF BAND MATRICES NM,9-67(285)
 F2 EIVALUES OF SYMM. TRIDIAG. MATRIX NM,9-67(388)
 F2 EIVALUES-EIVECTORS OF REAL MTRX NM,11-68(3)
 F2 SYM EIEROBLEM $A^*X=LAM*B^*X$ NM,11-68(102)
 F2 RATIONAL QR FOR SYM TRIDIAG NM,11-68(268)
 F2 MOD. LR.-CMPLX HESSENBERG MATRIX NM,12-68(372)
 F2 IMPLICIT QL ALGORITHM NM,12-68(379)
 F2 BALANCING A MATRIX NM,13-69(293)
 F2 QR FOR REAL HESSENBERG MTRX NM,14-70(219)
 F2 EIGENVEC OF REAL AND COMP MAT NM,16-70(181)
 F2 SKEW-SYMMETRIC MATRIX NM,17-71(189)
 F2 JACOBI TYPE METHOD FOR COMPLEX SYM. (A) NM,25-76(347)
 F2 496 LZ ALG. FOR GEN. COMPLEX EIGENVALUES (F) TO,1-75(271),2-76(396)
 F2 506 EIGENVALUES OF HESSENBERG MATRIX (F) TO,2-76(275)
 F2 517 COND. NOS. OF EIGENVALUES (F) TO,3-77(186)
 F2 530 SKEW-SYMMETRIC MATRICES (F) TO,4-78(286)
 F2 535 GEN. EIGENPROBLEM FOR COMPLEX MAT. (F) TO,4-78()
 F2 538 EIGENSYSTEMS OF GEN. SYMM. MATRICES (F) TO,5-79(118)
 F2 560 JORDAN NORMAL FORM OF COMPLEX MATRIX (F) TO,6-80(437)
 F2 19 EIGENVALUES OF REAL MATRIX (A) ZH,13-72(283)
 F2 225(37)NO.18
 F2 225(37)NO.21
- F3 DETERMINANTS
 F3 82 DETERMINANT OF ORTHOGONAL MATRIX (F) AS,24-75(150)
 F3 41 DETERMINANT EVALUATION CA,4-61(176),6-63(520),
 F3 41 7-64(144),9-66(686)
 F3 159 DETERMINANT EVALUATION CA,6-63(104),6-63(739)
 F3 170 DETERMINANT-POLYNOMIAL ELEMENTS CA,6-63(165),6-63(450),
 F3 170 7-64(421)
 F3 224 EVALUATION OF DETERMINANT CA,7-64(243),7-64(702)
 F3 269 DETERMINANT BY GAUSSIAN ELM. CA,8-65(668),9-66(686)
- F4 SIMULTANEOUS LINEAR EQUATIONS
 F4 108 MULT. LINEAR REGRESSION IN LI SENSE (F) AS,26-77(106),27-78(378)
 F4 108 MULT. LINEAR REGRESSION IN LI SENSE (F) AS,26-77(106)
 F4 132 LEAST ABSOLUTE VALUE ESTIMATION (F) AS,27-78(363)
 F4 4 GAUSSIAN ELIMINATION BT,2-62(256),3-63(60)
 F4 7 LINEAR SYSTEM WITH BAND MATRIX BT,3-63(207)
 F4 12 GAUSSIAN ELIMINATION BT,5-65(64)
 F4 68 IT REFINEMENT-LEAST SQR SOLN BT,8-68(20)
 F4 27 LIN EQ WITH SYM NON-DEF MATRIX BT,10-70(386)
 F4 16 CROUT WITH PIVOTING CA,3-60(507),3-60(540),
 F4 16 4-61(154)
 F4 17 SOLVE TRIDIAGONAL MATRIX CA,3-60(508)
 F4 24 SOLVE TRIDIAGONAL MATRIX CA,3-60(602)
 F4 43 CROUT WITH PIVOTING CA,4-61(176),4-61(182),
 F4 43 6-63(445)
 F4 92 SIMULT. EQNS.-ITERATIVE SOLN. CA,5-62(286)
 F4 107 GAUSSIAN ELIMINATION CA,5-62(388),6-63(39),
 F4 107 6-63(445)
 F4 126 GAUSSIAN ELIMINATION CA,5-62(511)
 F4 135 CROUT WITH EQUILIBRATION CA,5-62(553),5-62(557),
 F4 135 7-64(421),8-65(104)
 F4 195 BAND SOLVE CA,6-63(441),15-72(1074)
 F4 220 GAUSS-SEIDEL CA,6-63(739),7-64(349)
 F4 238 CONJUGATE GRADIENT METHOD CA,7-64(481)
 F4 288 LINEAR DIOPHANTINE EQUATIONS CA,9-66(514)
 F4 290 EXACT SOLUTION OF LINEAR EQNS. CA,9-66(683)
 F4 328 CHEBY SOLN-OVERDET LINEAR SYS CA,11-68(428),12-69(326)
 F4 358 SING.VAL.DECOMP.-COMPLEX MATRIX CA,12-69(564)
 F4 406 SOLUTION BY RESIDUE ARITH. (F) CA,14-71(180),16-73(311)
- F4 408 SPARSE MAT. PACKAGE(PART 1) (F) CA,14-71(265),16-73(311)
 F4 408 16-73(578)
 F4 408 A SPARSE MATRIX PACKAGE (PART 1) (F) CA,14-71(265),16-73(311)
 F4 408 A SPARSE MATRIX PACKAGE (F) CA,14-71(265),16-73(311)
 F4 408 16-73(578),3-77(303)
 F4 16-73(578),3-77(303)
 F4 4-78(295),6-80(456)
 F4 408 SPARSE MAT. PACKAGE (PART 1) (F) CA,14-71(265),16-73(311)
 F4 16-73(578),3-77(303),
 F4 4-78(295)
 F4 423 LINEAR EQUATION SOLVER CA,15-72(274)
 F4 432 SOLUTION TO EQUATION $AX+XB=C$ CA,15-72(820)
 F4 449 SOL LIN PROG PROB 0-1 VAR (F) CA,16-73(445)
 F4 470 ALMOST TRIDIAGONAL MATRIX (F) CA,16-73(760)
 F4 478 OVERDETERMINED SYSTEM IN LI NORM (F) CA,17-74(319),18-75(277)
 F4 DETERMINATION OF SSOR-SI ITER. PARAM.(F) CC,10-75(194)
 F4 SOLUTION WITH REL ERR ESTIMATE CJ,11-68(92)
 F4 51 GENL. 3-TERM LINEAR SYSTEMS CJ,13-70(324)
 F4 69 SYMMETRIC THREE TERM SYSTEMS CJ,14-71(444)
 F4 TRIDIAGONAL SYSTEMS (A) CJ,15-72(356)
 F4 80 QUINDIAGONAL SYSTEMS (A) CJ,16-73(278)
 F4 82 PARTIAL PIVOTING, REDUCED STORAGE (A) CJ,17-74(377)
 F4 84 TRIDIAGONAL SYSTEMS COUPLED ALGORITHM(A) CJ,17-74(378)
 F4 84 COUPLED ALG. FOR TRIDIAGONAL SYSTEMS (A) CJ,17-74(378),20-77(91)
 F4 SOL. OF CERTAIN BANDED LINEAR SYSTEMS(F) CJ,19-76(184)
 F4 BANDED LINEAR SYSTEMS (F) CJ,19-76(184)
 F4 LARGE SPARSE LINEAR EQUATIONS (P) CJ,19-76(352)
 F4 SOLUTION OF TRIDIAGONAL LINEAR SYS. (A) CJ,20-77(181)
 F4 96 LIN. EQNS. FROM GALERKIN METHOD (A) CJ,20-77(371)
 F4 98 TRIDIAGONAL SYSTEMS OF EQUATIONS (A) CJ,20-77(376)
 F4 108 TRIDIAGONAL LINEAR SYSTEMS (A) CJ,22-79(283)
 F4 ERROR BOUNDS FOR SYS LIN EQ CP,6-70(28)
 F4 ANTISYMMETRICAL NETWORKS (A) CP,11-73(45)
 F4 LIN CONSTRAINED OPTIM. PROBS(F) HA,AERE-R6354
 F4 SOL LARGE SPARSE SETS LIN EQ (F) HA,AERE-R6545
 F4 SOL OF SPARSE SETS LIN EQ (F) HA,AERE-R6844
 F4 SOL OF LINEAR EQUATIONS (F) HA,AERE-R6899
 F4 SPARSE, SYMM, & POS-DEFINITE(F) HA,AERE-R7119
 F4 TRIDIAGONAL SYSTEMS (A) JC,20-73(27)
 F4 VANDERMONDE SYSTEM OF EQUATION MC,24-70(893)
 F4 CONJUGATE GRADIENT METHOD NM,5-63(195)
 F4 LEAST SQUARES SOLUTION NM,7-65(271)
 F4 ELM.WITH WEIGHTED ROW COMB. NM,7-65(341)
 F4 ITER.REFIN.-SOLN.OF POS.DEF.MTX NM,8-66(206)
 F4 REAL AND COMPLEX LINEAR SYSTEM NM,8-66(222)
 F4 SYMM.AND UNSYMM.BAND EQUATIONS NM,9-67(285)
 F4 VANDERMONDE SYSTEMS NM,18-71(44)
 F4 DECOMPOSITION OF A SYMMETRIC MATRIX (A) NM,27-76(95)
 F4 495 OVERDETERMINED SYST IN CHEBYSHEV NORM(F) TO,1-75(264)
 F4 512 POS. DEF. SYM. QUINDIAGONAL LIN. (F) TO,3-77(96)
 F4 522 EXACT SOL. OF INTEGER LINEAR SYSTEMS (F) TO,3-77(404)
 F4 533 SPARSE GAUSSIAN ELIMINATION (F) TO,4-78()
 F4 544 WEIGHTED LEAST SQ. ITER. REFINEMENT (F) TO,5-79()
 F4 546 ALMOST BLOCK DIAGONAL SYSTEMS (F) TO,6-80(88)
 F4 551 LI SOLUTION OF OVERDETERMINED SYSTEMS(F) TO,6-80(228)
 F4 552 CONSTRAINED LI LINEAR APPROXIMATION (F) TO,6-80(231)
 F4 563 LINEAR CONSTRAINED DISCR. LI APPROX. (F) TO,6-80()
 F4 564 TEST PROBLEM GEN. FOR LI APPROX. (F) TO,6-80()
 F4 29 SOKOLOV'S METHOD (A) ZM,14-74(127)
 F4 37 SOL. OF SPARSE LINEAR EQ. SYSTEMS (A) ZM,14-74(841)
- F5 ORTHOGONALIZATION
 F5 10 USE OF ORTHOGONAL POLYNOMIALS AS,17-68(283),20-71(118)
 F5 10 20-71(216)
 F5 42 ORTHOGONAL POLYNOMIALS AS,20-71(209)
 F5 46 GRAM-SCHMIDT ORTHOGONALIZATION AS,20-71(335)
 F5 127 ORTHONORMALIZATION CA,5-62(511),13-70(122)
 F5 358 SING.VAL.DECOMP.-COMPLEX MATRIX CA,12-69(564)
 F5 SCHMIDT ORTHONORMALIZATION CP,1-66(159)
- G1 SIMPLE CALCULATIONS ON STATISTICAL DATA
 G1 16 ESTIMATION FROM NORMAL DATA AS,18-69(110)
 G1 16 ESTIMATION FROM NORMAL DATA AS,18-69(110),26-77(122)
 G1 19 VARIANCE OF FACTORIAL TABLE AS,18-69(199)
 G1 22 THE INTERACTION ALGORITHM AS,18-69(283)
 G1 23 CALCULATION OF EFFECTS AS,18-69(287)
 G1 25 MEANS FROM VARIANCE ANALYSIS AS,18-69(294)
 G1 29 THE RUNS UP AND DOWN TEST AS,19-70(190),25-76(193)
 G1 31 BINOMIAL SEQUENTIAL SAMPLING AS,19-70(197)
 G1 33 HYPERGEOMETRIC SAMPLE SIZES AS,19-70(287)
 G1 35 FINITE POPULATION PROBAB. (A) AS,20-71(99),20-71(346)
 G1 35 21-72(352)
 G1 35 PROB. FROM FINITE POPULATIONS (A) AS,20-71(99),20-71(346)
 G1 36 EXACT CONFIDENCE LIMITS AS,20-71(105)
 G1 41 MATRIX MEAN AND DISPERSION AS,20-71(206)
 G1 48 UNCERTAINTY FUN. FOR BINARY SEQ AS,21-72(97)
 G1 53 WISHART VARIATE GENERATOR (F) AS,21-72(341)
 G1 55 MANN-WHITNEY U-STATISTIC (P) AS,21-72(348)
 G1 59 HYPERGEOMETRIC PROBABILITIES(F) AS,22-73(130)
 G1 62 MANN-WHITNEY U-STATISTIC (F) AS,22-73(269)
 G1 67 ABSORPTION IN SEQ BINOM SAMP(F) AS,23-74(83)
 G1 68 TRUNCATED NEG. BIN. DISTR. (F) AS,23-74(87)
 G1 70 PERCENTAGE PTS OF NORM DIST (F) AS,23-74(96)
 G1 72 MEAN VECTORS & DISPERSION MATRICES (F) AS,23-74(234)
 G1 77 NULL DIST OF LARGEST ROOT STATISTIC (F) AS,23-74(458)
 G1 78 THE MEDIANCENTRE (F) AS,23-74(466),24-75(390)
 G1 80 SPHERICAL STATISTICS (A) AS,24-75(144)
 G1 81 CIRCULAR STATISTICS (A) AS,24-75(147)
 G1 84 MULTIVARIATE SKEWNESS AND KURTOSIS (F) AS,24-75(262)

- G1 85 CRITICAL VALUES OF SIGN TEST (A) AS,24-75(265)
G1 86 VON MISES DISTRIBUTION FUNCTION (A) AS,24-75(268)
G1 89 UPPER TAIL PROB. OF SPEARMAN'S RHO (F) AS,24-75(377)
G1 91 PERCENTAGE PTS. OF CHI SQUARED DIST. (F) AS,24-75(385)
G1 92 SAMPLE SIZE FOR DIST.-FREE TOL. INT. (F) AS,24-75(388)
G1 93 NULL DIST. OF ANSARI-BRADLEY W STAT. (F) AS,25-76(75)
G1 94 COEFFICIENTS OF THE ZONAL POLYNOMIALS (F) AS,25-76(82)
G1 95 EST. LOC. & SCALE PARAM. GROUPEDED DATA (F) AS,25-76(88)
G1 99 FITTING JOHNSON CURVES BY MOMENTS (F) AS,25-76(180)
G1 100 NORMAL-JOHNSON & JOHNSON-NORMAL TRANS. (F) AS,25-76(190)
G1 101 DISTRIBUTION-FREE CONFIDENCE INTER. (F) AS,25-76(309)
G1 107 CHAR. FOR SEQUENTIAL SAMPLING PLANS (F) AS,26-77(98)
G1 111 PERCENTAGE POINTS OF THE NORMAL DIST. (F) AS,26-77(118)
G1 112 EXACT DIST. FROM TWO-WAY TABLES (F) AS,26-77(199)
G1 112 EXACT DIST. FROM TWO-WAY TABLES (F) AS,26-77(199), 27-78(109)
G1 114 NUMERATOR OF ORDINAL MEAS. OF ASSOC. (F) AS,26-77(211)
G1 115 EXACT 2-SIDED CONFID. LIMITS 2X2 TAB. (F) AS,26-77(214)
G1 116 TETRACHORIC CORRELATION (F) AS,26-77(343)
G1 118 APPROXIMATE RANKITS (F) AS,26-77(362)
G1 119 JOINT FREQUENCY DISTRIBUTIONS (F) AS,26-77(364)
G1 122 WEIGHTS FOR 1-SIDED MULTIVAR. INF. (F) AS,27-78(100)
G1 124 SAMPLE SIZES FOR DIST-FREE TOL. LIMIT (A) AS,27-78(188)
G1 125 MAX. LIKELIHOOD EST. EXPON. SURVIVAL (F) AS,27-78(190)
G1 126 PROBABILITY INTEGRAL OF NORMAL RANGE (F) AS,27-78(197)
G1 128 APPROX. COVAR. MATRIX OF NORMAL ORDER (F) AS,27-78(206)
G1 129 POWER FUNCT., COMPARING BINOM. DIST. (F) AS,27-78(212), 29-80(118)
G1 130 MOVING STATISTIC FOR SCATTER PLOTS (F) AS,27-78(354)
G1 131 FREQ. DIST., STRUCTURED CODE SETS (F) AS,27-78(359)
G1 134 GEN. BETA RANDOM VARIABLE (F) AS,28-79(90)
G1 136 A K-MEANS CLUSTERING ALGORITHM (F) AS,28-79(100)
G1 137 DEPEND. SAMPLES MULTIVAR. DENSITY (F) AS,28-79(109)
G1 138 MAX LIKELIHOOD EST. CONF. NORMAL DATA (F) AS,28-79(185)
G1 139 MAX LIKELIHOOD EST. LINEAR MODEL (F) AS,28-79(214), 29-80(228)
G1 146 JOINT PROB. OF SELECTION, PPS SAMPLE (F) AS,29-80(107)
G1 148 THE JACKKNIFE (F) AS,29-80(115)
G1 149 AMALGAMATION OF MEANS, SIMPLE ORDER (F) AS,29-80(209)
G1 150 SPECTRUM EST. FOR A COUNTING PROCESS (F) AS,29-80(211)
G1 151 SPECTRAL EST. FOR BIVARIATE COUNTING (F) AS,29-80(214)
G1 152 CUMULATIVE HYPERGEOMETRIC PROB. (F) AS,29-80(221)
G1 153 TAIL PROB. OF DURBIN-WATSON STATISTIC (A) AS,29-80(224)
G1 208 DISCRETE CONVOLUTION CA,6-63(615)
G1 212 DETERMINE DISTRIB. FCN. FROM DATA CA,6-63(617)
G1 289 CONFIDENCE INTERVAL FOR A RATIO CA,9-66(514)
G1 330 FACTORIAL ANALYSIS OF VARIANCE CA,11-68(431)
G1 359 FACTORIAL ANALYSIS OF VARIANCE CA,12-69(631), 13-70(449)
G1 451 CHI-SQUARE QUANTILES (F) CA,16-73(483), 18-75(116)
G1 TAIL AREA PROB. FOR 2X2 TABLE CB,9-65(56)
G1 FITTING THEORETICAL EXPRESSION CC,3-72(296)
G1 KAC AND RENEYI TESTS (F) CC,12-76(173)
G1 CJ,9-66(212), 9-67(416)
G1 CJ,17-74(378)
G1 CJ,21-78(270)
G1 CP,16-76(379)
G1 TO,3-77(183)
G1 21-72(352), 26-77(221)
G1 ZM,12-72(499)
G1 ZM,14-74(335)
G1 ZM,14-74(343)
G1 ZM,14-74(471)
G1 ZM,14-75(639)
G1 ZM,14-75(655)
G2 CORRELATION AND REGRESSION ANALYSIS
G2 17 RECIPROCAL OF MILLS RATIO AS,18-69(115)
G2 38 BEST SUBSET SEARCH AS,20-71(112)
G2 49 AUTOCORRELATION FUN FOR BIN SEQ AS,21-72(100)
G2 54 KENDALL'S S FREQUENCY DIST. (F) AS,21-72(345)
G2 58 EUCLIDIAN CLUSTER ANALYSIS (F) AS,22-73(126), 24-75(160)
G2 69 SPACE-TIME CLUST. IN EPIDEM. (F) AS,23-74(92)
G2 71 UPPER TAIL PROB KENDALLS TAU (F) AS,23-74(98)
G2 75 LINEAR LEAST SQUARES REGRESSION (A) AS,23-74(448)
G2 75 LINEAR LEAST SQUARES REGRESSION (A) AS,23-74(448), 25-76(323)
G2 79 GRAM-SCHMIDT REGRESSION (A) AS,23-74(470)
G2 87 EST. OF CORR. IN CONTINGENCY TABLES (F) AS,24-75(272)
G2 87 POLYCHORIC EST. OF CORR. IN CONT. TAB. (F) AS,24-75(272), 26-77(121)
G2 90 ONE-SIDED MULTI-VARIABLE INFERENCE (F) AS,24-75(380)
G2 104 BLUS RESIDUALS (A) AS,25-76(317)
G2 105 COVARIANCE SELECTION MODEL TO A MATRIX (F) AS,26-77(88)
G2 120 TWO-WAY UNBALANCED MANOVA (F) AS,27-78(92)
G2 135 MIN-MAX EST. LIN. MULT. REGR. (F) AS,28-79(93)
G2 139 MAX LIKELIHOOD EST. LINEAR MODEL (F) AS,28-79(195)
G2 141 INVERSE OF SYMM. MATR. IN REGR. MOD. (F) AS,28-79(214)
G2 39 CORRELATION COEFFICIENTS CA,4-61(152)
G2 142 TRIANGULAR REGRESSION CA,5-62(603)
G2 366 REGRESSION-DIR. PROD. OF MATRICES CA,12-69(687)
G2 367 ANALYSIS OF VAR.-DIRECT EXPMT. CA,12-69(688)
G2 434 RxC CONTINGENCY TABLES (F) CA,15-72(991), 17-74(326)
G2 434 18-75(117), 2-76(108)
G2 NON-LINEAR LEAST SQUARES CC,2-71(77), 5-73(308)
G2 SINGLE-LINK CLUSTER METHOD (F) CJ,16-73(30)
G2 76 HIERARCHICAL CLUSTERING (F) CJ,16-73(93)
G2 FREQ. DIST. OF RANK CORRELATION CP,8-71()
G2 CLUST. 1-DIMEN. ORDERED DATA (A) CP,11-73(175)
G2 SAMP PARAM FOR ORTH LIN REG (F) CP,13-74(17)
G2 39 CLUSTERWISE LINEAR REGRESSION (F) CP,22-79(367)
G2 39 CORRECTED SUMS OF SQS. FOR ANAL. VAR. (A) ZM,14-75(647)
G2 41 INTERDEP. EXAMS. BY ANAL. REGRESSION (A) ZM,15-76(125)
G2 44 MAHALANOBIS DIST. OR RES. SUM SQS. (A) ZM,15-76(215)
G5 RANDOM NUMBER GENERATORS
G5 98 SPECTRAL TEST FOR CONG. RAND. NO. GEN. (F) AS,25-76(173)
G5 98 SPECTRAL TEST FOR CONG. RAND. NO. GEN. (F) AS,25-76(173), 27-78(375)
G5 98 SPECTRAL TEST FOR CONG. RAND. NO. GEN. (F) AS,25-76(173), 25-76(324)
G5 106 DISTRIBUTION OF NON-NEG QUAD. FORMS (F) AS,26-77(92)
G5 121 RANDOM NORMAL CA,5-62(482), 8-65(556)
G5 133 RANDOM FLAT CA,5-62(553), 5-62(606), 6-63(105), 6-63(167)
G5 133 CA,6-63(444), 8-65(556)
G5 200 RANDOM NORMAL CA,7-64(701)
G5 247 QUASI-RANDOM POINT SEQUENCE CA,8-65(605), 9-66(687), 15-72(1072)
G5 266 PSEUDO-RANDOM NUMBERS CA,8-65(606)
G5 266 CA,10-67(40)
G5 267 RANDOM NORMAL DEVIATES CA,11-68(498), 12-69(281)
G5 294 UNIFORM RANDOM CA,11-68(819)
G5 334 NORMAL RANDOM DEVIATES CA,13-70(49)
G5 342 POISSON RANDOM NUMBERS CA,13-70(49), 15-72(467)
G5 369 POISSON RANDOM NUMBERS CA,13-70(326), 15-72(468)
G5 370 RANDOM NUMBER GENERATOR CA,15-72(355), 17-74(325)
G5 381 RANDOM VECTOR IN SOLID ANGLE CA,16-73(51)
G5 425 RANDOM CORRELATED NORM VAR (F) CA,17-74(703)
G5 441 DEVIATES FROM DIPOLE DIST. (F) CB,9-65(105)
G5 488 GAUSSIAN PSEUDO-RANDOM (F) CC,9-75(360)
G5 RANDOM UNIFORM CC,12-76(163)
G5 GENERAL PURPOSE MONTE-CARLO PROGRAM (F) CJ,6-63(279)
G5 SEQUENTIAL RANDOM INTEGER GENERATOR (F) CP,9-72(63)
G5 RANDOM SAMPLES, VARIOUS DISTRIB. JC,18-71(381)
G5 PSEUDO-RANDOM NUMBER GENERATION JC,20-73(461)
G5 TALSNOTHE PSD-RAN NUM GEN JC,20-73(469)
G5 PSEUDORANDOM NUMBER ALGOR. (F,M)
G5 ASYMPTOT. RNDM TAUSWORTH SEQ(F)
C6 PERMUTATIONS AND COMBINATIONS
C6 56 PERMUTATIONAL SIGNIFICANCE (A) AS,22-73(112)
C6 88 ALL COMB. OF N THINGS R AT A TIME (F) AS,24-75(374)
C6 PARTITION FCNS. (MODULO D) BT,9-69(83)
C6 IN PLACE PERMUTATION (A) BT,12-72(318)
C6 71 PERMUTATIONS CA,4-61(497), 5-62(209), 5-62(439)
C6 71 CA,5-62(208), 5-62(209), 5-62(440)
C6 86 PERMUTATIONS CA,5-62(209), 5-62(440), 5-62(514), 10-67(452)
C6 86 CA,5-62(344), 5-62(557), 5-62(606)
C6 87 PERMUTATION GENERATOR CA,5-62(346), 5-62(514), 10-67(452)
C6 87 CA,5-62(434), 5-62(514), 5-62(606)
C6 94 COMBINATIONS CA,5-62(551), 10-67(452)
C6 94 CA,6-63(68), 6-63(385)
C6 102 PERMUTATIONS IN LEXIC. ORDER CA,6-63(103), 6-63(449)
C6 102 CA,6-63(103), 6-63(449)
C6 102 CA,6-63(103), 6-63(450)
C6 102 COMB. OF M THINGS N AT A TIME CA,6-63(161), 6-63(450), 6-63(618)
C6 161 COMBS. 1,2,UP TO N AT A TIME CA,6-63(161), 6-63(450), 6-63(619)
C6 161 CA,6-63(517), 8-65(556), 10-67(452)
C6 202 PERMUTATIONS IN LEXIC. ORDER CA,7-64(420), 8-65(445)
C6 202 CA,7-64(585)
C6 235 RANDOM PERMUTATION CA,8-65(104), 8-65(670)
C6 242 PERMUTATIONS WITH REPETITIONS CA,10-67(450)
C6 250 INVERSE PERMUTATION CA,10-67(452), 12-69(638)
C6 306 PERMUTATIONS WITH REPETITIONS CA,10-67(729)
C6 308 PERMUT. IN PSEUDOLEXIC. ORDER CA,11-68(117), 16-73(577)
C6 317 PERMUTATION CA,11-68(430), 12-69(187)
C6 323 PERM. IN LEXICOGRAPHIC ORDER (A) CA,12-69(634), 13-70(376)
C6 329 DISTR OF INDISTINGUISHABLE OBJ CA,12-69(634)
C6 361 PERMANENT FNC. OF SQUARE MATRIX CA,12-69(634)
C6 362 RANDOM PERMUTATIONS CA,13-70(368), 13-70(376)
C6 382 COMBINATIONS OF M OUT OF N CA,13-70(368), 13-70(376)
C6 383 PERMUTATIONS WITH REPETITIONS CA,16-73(485)
C6 452 ENUMER. COMB. M OUT OF N OBJ (F) CA,16-73(690)
C6 466 FOUR COMBINATORIAL ALGOR. (P) CB,9-65(104)
C6 ALL PERMUTATIONS OF N OBJECTS CC,1-70(415)
C6 ALG GENERATING PERMUTATION CJ,10-67(311)
C6 28 PERMUTNS OF VECTORS-LEXIC ORDER CJ,10-67(311)
C6 29 PERMUTNS OF VECTORS CJ,10-67(311)
C6 30 FAST PERMUTN OF VECTORS CJ,14-71(136)
C6 PERMUTATION SEQUENCE, PART 2 CJ,16-73(57)
C6 OCCUPANCY OF RECTANG. ARRAY (A) CJ,19-76(156)
C6 GENERATING PERMUTATIONS (A) CJ,19-76(156)
C6 GENERATING PERMUTATIONS (A) CJ,21-78(373)
C6 104 GEN. OF ALPHABET. ORDERED SEQUENCES (A) CP,12-74(269)
C6 19 STEP-CYCLE GENERATION (A) MC,24-70(955)
C6 SOLID PARTITIONS TO,3-77(180)
C6 515 GEN. OF VECTOR FROM LEXICO. INDEX (F)
G7 SUBSET GENERATORS
G7 15 SINGLE LINKAGE CLUSTER ANALY AS,18-69(106)
G7 113 NON-HIERARCHICAL CLASSIFICATION (F) AS,26-77(206)
G7 81 SUBSEQUENCES CA,5-62(166)
G7 82 SUBSEQUENCES CA,5-62(167)
G7 477 SET-PARTITIONS TO R SUBSETS (A) CA,17-74(224)
G7 482 TRANSIVITY SETS (A) CA,17-74(470)
G7 47 A CLUSTERING ALGORITHM CJ,13-70(113), 13-70(219)
G7 47 13-70(326)
G7 49 INDEXING SUBARRAYS CJ,13-70(208), 13-70(219)
G7 52 HIERARCHICAL CLUSTERING ALGORITHM CJ,13-70(325)

G7	65 IMPROVED CLUSTERING ALGORITHM	CJ,14-71(104)	H	20 K-TH BEST ROUTE IN GRAPH (A)	CP,12-74(273)
G7	29 DOUBLY & MULTIPLY RES. PARTITIONS (A)	CP,16-76(163)	H	23 AN IMPROVED MATCHING ALGORITHM (P)	CP,13-74(389)
H	OPERATIONS RESEARCH, GRAPH STRUCTURES				
H	13 MINIMUM SPANNING TREE	AS,18-69(103)	H	24 DERIVING CHROMATIC POLYNOM. OF GRAPH (F)	CP,14-75(195),16-76(291)
H	14 PRINTING MINIMUM SPANNING TREE	AS,18-69(105)	H	25 MIXED INTEGER CONVEX PROGRAMMING (A)	CP,14-75(389)
H	40 MINIMAL SPANNING TREE	AS,20-71(204)	H	26 TRANSITIVE CLOSURE OF A DIGRAPH	CP,15-75(75)
H	102 ULTRAMETRIC DISTS. FOR S.L. DENDRO. (F)	AS,25-76(313)	H	OPTIMIZATION FOR SPECIAL NETWORKS (F)	CP,16-76(77)
H	140 CLUSTERING OF NODES OF DIR. GRAPHS (F)	AS,28-79(206)	H	27 K-TH BEST ROUTE PROBLEM (A)	CP,16-76(139)
H	10 LINEAR PROGRAMMING	BT,4-64(189),6-66(82),	H	28 GENERATING GRAPHS OF GIVEN PARTITION (F)	CP,16-76(153)
H	10	12-72(267)	H	31 NO. OF PATHS IN A NETWORK	CP,17-76(93)
H	TRANSITIVE CLOSURE ALGORITHM	BT,10-70(76)	H	31 NUMBER OF PATHS IN A NETWORK (A)	CP,17-76(93)
H	STABLE MARRIAGE ASSIGNMENT	BT,10-70(295)	H	37 0-1 SINGLE KNAPSACK PROBLEM (F)	CP,21-78(81)
H	DERIVATION OF FRISCH ALGORITHM	BT,11-71(94)	H	41 TRAVERSAL OF UNDIRECTED GRAPH (F)	CP,24-80(269)
H	PARTITION OF UNDIRECTED GRAPH	BT,12-72(161)	H	PARSING OF GRAPH-REP PICTURES	JC,17-70(453)
H	UNIQUE INCID. MAT. OF GRAPH (A)	BT,14-74(209)	H	EDGE AND LINE OPERATOR (A)	JC,20-73(634),21-74(350)
H	OPTIMUM ASSIGNMENT (A)	BT,19-79(289)	H	MIN LENGTH FULL STEINER TREES	MC,23-69(521)
H	27 ASSIGNMENT PROBLEM	CA,3-60(603),6-63(618),	H	SIGN STABILITY AND Z-MAXIMUM MATCHINGS(A)NM	NM,28-77(273)
H	27	6-63(739)	H	520 REVISED SIMPLEX METHOD (F)	TO,3-77(295)
H	40 CRITICAL PATH SCHEDULING	CA,4-61(152),4-61(392),	H	548 SOLUTION OF THE ASSIGNMENT PROBLEM (F)	TO,6-80(104)
H	40	5-62(513),7-64(349)	H	557 LINEAR GOAL PROGRAMMING PROBLEMS (F)	TO,6-80(429)
H	69 CHAIN TRACING	CA,4-61(392)	H	558 MULTIFACILITY LOCATION, RECTIL. DIST.(F)	TO,6-80(450)
H	83 CLASSIFICATIONS	CA,5-62(167)	H	562 SHORTEST PATH LENGTHS (F)	TO,6-80(450)
H	96 ANCESTOR	CA,5-62(344),6-63(104)	H	3 ZERO-ONE INTEGER LINEAR PROG.	ZM,11-69(111),11-70(513)
H	97 SHORTEST PATH	CA,5-62(345)	H	5 TRANSPORTATION PROBLEM	ZM,11-70(221),12-71(131)
H	119 PERT NETWORK	CA,5-62(436),8-65(330)	H	9 TIME TRANSPORTATION PROBLEM	ZM,12-71(347)
H	141 FIND PATH	CA,5-62(556)	H	14 ZERO-ONE INTEGER LINEAR PROG.	ZM,13-72(133)
H	153 INTEGER PROGRAMMING	CA,6-63(68),6-63(449)	H	16 0-1 INTEGER LIN. PROC. (A)	ZM,13-72(133)
H	217 MIN. EXCESS COST CURVE	CA,6-63(737),11-68(573)	H	17 SHORTEST PATH FRM FIXED NODE(A)	ZM,13-72(275)
H	219 TOPOLOGICAL ORDERING	CA,6-63(738)	H	18 SHORTEST PATH FRM FIXED NODE(A)	ZM,13-72(279)
H	219 TOPOLOGICAL ORDERING FOR PERT NETS. (A)	CA,6-63(738),3-77(303)	H	20 CYCLES OF A GRAPH (A)	ZM,13-73(399)
H	248 NETFLOW	CA,8-65(103),11-68(633)	H	23 DISCRETE OPTIMIZATION (A)	ZM,13-73(411)
H	248	11-68(633)	H	24 ONE-DIM. KNAPSACK FUNCTION (A)	ZM,13-73(415)
H	258 TRANSPORT	CA,8-65(381),8-65(445),	H	25 DISCRETE OPTIMIZATION (A)	ZM,13-73(541)
H	258	10-67(453)	H	26 QUEING SYSTEM (A)	ZM,13-73(548)
H	263 INTEGER PROGRAMMING-GOMORY1	CA,8-65(601),13-70(326)	H	36 TRANSITIVE CLOSURE OF A GRAPH (A)	ZM,14-74(477)
H	285 MUTUAL PRIMAL-DUAL METHOD	CA,9-66(326),10-67(453)	H	42 ENCLOSURE OF PT. TO MIN. SPAN. TREE (A)	ZM,15-76(135)
H	286 EXAMINATION SCHEDULING	CA,9-66(433),9-66(795)	H	45 TRAVELING-SALESMAN PROBLEM (A)	ZM,15-76(223)
H	293 TRANSPORTATION PROBLEM	CA,9-66(869),10-67(453)	I5	INPUT - COMPOSITE	
H	293	11-68(271)	I5	43 VARIABLE FORMAT IN FORTRAN	AS,20-71(213),20-71(346)
H	324 MAXFLOW (A)	CA,11-68(117),16-73(309)	I5	239 FREE-FIELD READ	CA,7-64(481)
H	333 MINIT ALGORITHM FOR LIN PROG(A)	CA,11-68(437),12-69(408)	I5	249 OUTREAL N	CA,8-65(104)
H	333	14-71(50),16-73(310)	I5	335 BASIC I/O PROCEDURES	CA,11-68(567)
H	336 NETFLOW	CA,11-68(631),13-70(192)	I5	CODNUM, CHANGE OF PUNCH CODE	CC,2-71(168)
H	341 LINEAR PGNS. IN 0-1 VARIABLES	CA,11-68(782),12-69(692)	I5		1962(236)
H	341	13-70(263)	J6	PLOTTING	
H	350 SIMPLEX METHOD-LU DECOMPOSITION	CA,12-69(275)	J6	21 SCALE SELECTION FOR COMPUTER PLOTS (F)	AS,18-69(206),20-71(118)
H	354 SPANNING TREE GENERATOR	CA,12-69(511)	J6	21	23-74(248)
H	360 SHORTEST PATH FOREST-TOPOL. ORD.	CA,12-69(632)	J6	30 HALF NORMAL PLOTTER	AS,19-70(192),20-71(118)
H	394 DECISION TABLE TRANSLATION	CA,13-70(571),15-72(107)	J6	30	21-72(351)
H	397 INTEGER PROGRAMMING PROBLEM	CA,13-70(620),15-72(469)	J6	44 SCATTER DIAGRAM PLOTTING (F)	AS,20-71(327),23-74(248)
H	399 SPANNING TREE	CA,13-70(621)	J6	45 HISTOGRAM PLOTTING (F)	AS,20-71(332),22-73(274)
H	411 STABLE MARRIAGE PROBLEM	CA,14-71(491)	J6	61 SIX LINE PLOTS (F)	AS,22-73(265)
H	415 ASSIGNMENT PROBLEM	CA,14-71(805)	J6	61 SIX-LINE PLOTS (F)	AS,22-73(265),26-77(368)
H	422 MINIMAL SPANNING TREE	CA,15-72(273)	J6	96 SCALING GRAPHS (F)	AS,25-76(94)
H	430 PREDOMINATORS IN DIRECTED GRAPH	CA,15-72(777)	J6	162 XY PLOTTER	CA,6-63(161),6-63(450),
H	431 QUAD. AND LINEAR PROGRAMMING(F)	CA,15-72(818),17-74(157)	J6	162	7-64(482)
H	431	17-74(590)	J6	278 GRAPH PLOTTER	CA,9-66(88)
H	447 GRAPH MANIPULATION (A)	CA,16-73(372)	J6	412 GRAPH PLOTTER	CA,14-71(492)
H	456 ROUTING PROBLEM (F)	CA,16-73(573),17-74(706)	J6	420 HIDDEN-LINE PLOTTING PROGRAM(F)	CA,15-72(100),16-73(578)
H	457 CLIQUES OF UNDIRECTED GRAPH(A)	CA,16-73(575)	J6	420	17-74(324),17-74(706)
H	459 ELEMENTARY CIRCUITS OF A GRAPH (A)	CA,16-73(632),18-75(119)	J6	463 DETERMINATION OF SCALES (F)	CA,16-73(639)
H	481 ARROW TO PREC. NETWORK TRANS(F)	CA,17-74(467)	J6	475 VISIBLE SURFACE PLOTTING (F)	CA,17-74(152),18-75(202)
H	491 BASIC CYCLE GENERATION (P)	CA,18-75(275)	J6	475	18-75(276),18-75(277),
H	492 GEN ALL CYCLES FROM SET OF BASIC CYC (P)	CA,18-75(310)	J6	475	1-75(381),2-76(109)
H		8-64(147),9-67(18)	J6	475	CA,17-74(153)
H	1 MINIMAL SPANNING TREE	CB,8-64(67),8-64(109),	J6	475 VISIBLE SURFACE PLOTTING (F)	CA,17-74(152),18-75(202)
H	OPTIMIZATION PROBLEMS	CC,3-72(159)	J6	475	18-75(276),18-75(277),
H	GENERATION OF CONNECTED GRAPHS (F)	CC,8-74(320)	J6	475	1-75(381),2-76(109),
H	14 PROCESSING EVENT NETWORK	CJ,9-66(323)	J6	475	5-79()
H	22 SHORTEST PATH-START TO END	CJ,10-67(306)	J6	483 MASKED 3-DIM PLOT WITH ROTATIONS (F)	CA,17-74(520),1-75(285)
H	23 SHORTEST PATH-START TO ANY	CJ,10-67(307)	J6	ISOMETRIC REP OF TWO-DIM MAT	CC,1-70(277)
H	24 NODES ON SHORTEST PATH	CJ,10-67(308)	J6	LAUE PHOTOGRAPHS	CC,1-70(293)
H	PROCESS CONTROL PROGRAMMING	CJ,13-70(70)	J6	PRINTER-PLOTTER ROUTINE	CC,2-71(55)
H	57 INDEX VALUE WITHIN A RANGE	CJ,13-70(425)	J6	PRINTER-PLOTTER ROUTINE	CC,2-71(470)
H	58 PRIMAL SIMPLEX LINEAR PROGRAM	CJ,13-70(426),14-71(215)	J6	PLOTTING PACKAGE (F)	CC,9-75(85)
H	59 PRIMAL SIMPLEX LINEAR PROGRAM	CJ,13-70(428),14-71(215)	J6	41 A CURVE PLOTTING PROCEDURES	CJ,12-69(291)
H	60 PRIMAL SIMPLEX LINEAR PROGRAM	CJ,13-70(429),14-71(215)	J6	54 APP OF STRAIGHT LINES	CJ,13-70(422)
H	61 PRIMAL SIMPLEX LINEAR PROGRAM	CJ,13-70(430),14-71(215)	J6	66 FRENCH CURVE PROCEDURE (F)	CJ,14-71(207),15-72(285)
H	DIGITAL SIMULATION (F,B)	CJ,16-73(118),16-73(382)	J6	75 CONTOUR PLOTTING (A)	CJ,15-72(382)
H	79 CAPACITATED TRANSPORT. PRBLM(A)	CJ,16-73(276)	J6	81 DENDROGRAM PLOTTING (F)	CJ,17-74(89),18-75(90)
H	DECOMPOSITION OF GOZINTO'S GRAPH (F)	CJ,17-74(245)	J6	DRAWING CONTOURS FROM ARB. DATA PTS (A)	CJ,17-74(318)
H	91 BALANCING A BINARY TREE (F)	CJ,19-76(360)	J6	DRAWING CONTOURS FROM ARB. DATA PTS. (A)	CJ,17-74(318),19-76(335)
H	99 BALANCED BINARY TREES (A)	CJ,20-77(378)	J6	100 VECTOR APPROXIMATION TO CURVES (F)	CJ,21-78(178)
H	INVOLUTIONIC AUTOMOR. OF GRAPH	CP,2-67(332)	J6	35 CONTOUR REPRESENTATIONS (F)	CP,19-78(381)
H	7 TRAVELING SALESMAN PROBLEM	CP,3-68(151),5-70(385)	J6	PLOTTING CUBIC SPLINE FN (F)	HA,AERE-R7470
H	8 BESTIMUNG ALLER MAXIMALEN	CP,3-68(240),4-69(75)	J6	531 CONTOUR PLOTTING (F)	TO,4-78(290)
H	9 PROC. TO MOD. ALG. OF MINTY, MOORE	CP,4-69(76)	K2	RELOCATION	
H	LOSUNG DES OPTIMUM-MIX-PROBLEME	CP,5-70(326)	K2	8 MAIN EFFECT OF MULTI-WAY TABLE	AS,17-68(277)
H	STRAY PROCESS	CP,6-70(139)	K2	9 CONSTRUCTION OF ADDITIVE TABLE	AS,17-68(279)
H	SPERNER SIMPLEX	CP,8-71(157)	K2	20 FORMATION OF TRIANGULAR ARRAY	AS,18-69(203)
H	SHORTEST PATH IN GRAPH	CP,8-71(171)	K2	28 TRANSPOSING MULTIWAY STRUCTURE	AS,19-70(115)
H	FLOW IN ANTISYMMETRICAL NETWORK	CP,8-71(191)	K2	173 TRANSFER ARRAY VALUES	CA,6-63(311),6-63(619)
H	MAXIMAL MATCHING IN GRAPHS (A)	CP,9-72(251)	K2	284 INTERCHANGE 2 BLOCKS OF DATA (A)	CA,9-66(326),2-76(392)
H	LINEAR ECONOMETRIC MODELS (F)	CP,10-72(33)	K2	302 TRANSPOSE VECTOR STORED ARRAY	CA,10-67(292),12-69(326)
H	DISTANCE COMP. IN GRAPHS (F)	CP,10-72(107)	K2	11 TRANSFORMATION OF AN OWN ARRAY	ZM,12-71(123)
H	STAND. FORMS PLANAR GRAPHS (A)	CP,10-72(121)			
H	ISOMORPHISM TO GRAPHS (P)	CP,11-73(159)			
H	NEG. CYCLES VALUED DIGRAPHS(A)	CP,11-73(169)			
H	MINIMAL COST FLOW PROBLEMS (A)	CP,11-73(275)			

L2	COMPILING		S13 20	4-61(182)
L2	39 ARRAY OF VARIABLE DIMENSION	AS, 20-71(115)	S13 108 EXPONENTIAL INTEGRAL	CA, 5-62(388), 5-62(393)
L2	65 INTERPRET STRUCTURE FORMULAE (F)	AS, 22-73(414)	S13 109 EXPONENTIAL INTEGRAL	CA, 5-62(388), 5-62(393)
L2	13 EVALUATION OF FCNAL EXPRESSION	BT, 5-65(133)	S13 385 EXPONENTIAL INTEGRAL	CA, 13-70(446), 13-70(448)
L2	TRANSFORMATION OF IDENTIFIERS	BT, 11-71(16)	S13 385	13-70(750), 15-72(1074)
L2	ACKERMANN FUNCTION	BT, 11-71(107)	S13 471 EXPONENTIAL INTEGRALS (A)	CA, 16-73(761)
L2	265 FIND PRECEDENCE FUNCTIONS	CA, 8-65(604)	S13 EXPONENTIAL INTEGRAL EXPANSION	CH, 6-XX(187)
L2	USE OF ALGOL 68 FOR TREES	CJ, 13-70(25)	S13 EI(X) BY CHEBYSHEV EXPANSION	NM, 4-63(413)
L2	ALG SEMANTICS FOR ALGOL 60	JC, 17-70(361)	S13 SIN INTEGRAL	NM, 9-67(381)
L2	CROSS COMPILER FOR POCKET CALC. (B)	CJ, 20-77(213)	S13 COS INTEGRAL	NM, 9-67(382)
			S13 556 EXPONENTIAL INTEGRALS (F)	TO, 6-80(420)
M1	SORTING			
M1	26 RANKING AN ARRAY OF NUMBERS (A)	AS, 19-70(111), 22-73(133)	S14 3 STUDENTS T-DISTRIBUTION	AS, 17-68(189), 18-69(118)
M1	23 SORT	CA, 3-60(601), 4-61(238)	S14 5 INTEG OF NON-CENTRAL T-DIST (F)	AS, 17-68(193), 18-69(118)
M1	63 SORT	CA, 4-61(321), 5-62(439),	S14 5	22-73(428)
M1	63	6-63(446)	S14 27 STUDENTS T-DISTRIBUTION	AS, 19-70(113)
M1	64 SORT	CA, 4-61(321), 5-62(439),	S14 32 INCOMPLETE GAMMA INTEGRAL	AS, 19-70(285)
M1	64	6-63(446)	S14 63 INCOMPLETE BETA INTEGRAL (F)	AS, 22-73(409)
M1	65 SORT	CA, 4-61(321), 5-62(439),	S14 63 INCOMPLETE BETA INTEGRAL (F)	AS, 22-73(409), 26-77(111)
M1	65	6-63(446)	S14 64 INVERSE INCOMP BETA FN RATIO(F)	AS, 22-73(411)
M1	76 SORT	CA, 5-62(48), 5-62(348)	S14 64 INVERSE INCOMP. BETA FUNC. RATIO (F)	AS, 22-73(411), 26-77(111)
M1	113 TREESORT	CA, 5-62(434)	S14 76 NON-CENTRAL T & BIVARIATE NORMAL PROB(F)	AS, 23-74(455)
M1	143 TREESORT	CA, 5-62(604)	S14 76 NON-CENTRAL T & BIVARIATE NORM PROB (F)	AS, 23-74(245), 27-78(379)
M1	144 TREESORT	CA, 5-62(604)	S14 76	28-79(113)
M1	151 LOCATE IN A LIST	CA, 6-63(68)	S14 76 NON-CENTRAL T & BIVARIATE NORM PROB (F)	AS, 23-74(245), 28-79(113)
M1	175 SHUTTLE SORT	CA, 6-63(312), 6-63(619)	S14 103 PSI (DIGAMMA) FUNCTION (F)	AS, 25-76(315)
M1	175	6-63(739), 7-64(296)	S14 109 INCOMPLETE BETA INT. & INVERSE ... (F)	AS, 26-77(111)
M1	201 SHELL SORT	CA, 6-63(445), 7-64(349),	S14 121 TRIGAMMA FUNCTIONS (F)	AS, 27-78(97)
M1	201	13-70(373)	S14 123 MIXTURES OF BETA DISTRIBUTIONS (F)	AS, 27-78(104)
M1	207 STRING SORT	CA, 6-63(615), 7-64(585)	S14 147 INCOMPLETE GAMMA INTEGRAL (F)	AS, 29-80(113)
M1	232 HEAPSORT	CA, 7-64(347)	S14 GAMMA FUNCTION	BT, 2-62(238)
M1	245 TREESORT 3	CA, 7-64(701), 8-65(445),	S14 31 GAMMA FUNCTION	CA, 4-61(105), 5-62(605)
M1	245	13-70(371)	S14 34 GAMMA FUNCTION	CA, 4-61(106), 5-62(391),
M1	271 QUICKSORT	CA, 8-65(669), 9-66(354)	S14 34	9-66(685)
M1	347 SORT WITH MINIMAL STORAGE (A AND F)	CA, 12-69(185), 13-70(54)	S14 54 GAMMA FUNCTION	CA, 4-61(180), 9-66(685)
M1	347	13-70(624), 2-76(290)	S14 80 GAMMA FUNCTION	CA, 5-62(166), 9-66(685)
M1	402 IMPROVED QUICKSORT (A)	CA, 13-70(693), 16-73(311)	S14 147 DERIVATIVE OF GAMMA FUNCTION	CA, 5-62(605), 6-63(168),
M1	410 PARTIAL SORTING	CA, 14-71(357)	S14 147	12-69(691)
M1	426 MERGE SORT ALGORITHM (A)	CA, 15-72(357), 17-74(706)	S14 179 BETA RATIO (F)	CA, 6-63(314), 10-67(375)
M1	426	2-76(290)	S14 179	17-74(156)
M1	TESTS OF QUICKSORT AND DESCENDANTS (F)	CA, 17-74(143), 3-77(204)	S14 221 GAMMA FUNCTION	CA, 7-64(143), 7-64(586),
M1	489 FINDING I-TH SMALLEST OF N ELEMENTS (A)	CA, 18-75(173), 2-76(301)	S14 221	9-66(685)
M1	SORTING OF INTEGERS	CB, 9-67(63)	S14 222 INCOMPLETE BETA FCN. RATIOS	CA, 7-64(143), 7-64(244)
M1	25 SORT BY RANKING ELEMENTS	CJ, 10-67(308), 12-69(408)	S14 225 GAMMA FCN WITH CONTROLLED ACCY.	CA, 7-64(295), 7-64(586)
M1	25	13-70(326)	S14 291 LOGARITHM OF GAMMA FCN.	CA, 9-66(684), 9-66(685),
M1	26 ORDER SUBSCRIPTS BY ELEMNT SIZE	CJ, 10-67(309), 12-69(408)	S14 291	11-68(14)
M1	26	13-70(326)	S14 309 GAMMA FCN.-ARBITRARY PRECISION	CA, 10-67(511)
M1	27 SORT ON PERMUTN OF SUBSCRIPTS	CJ, 10-67(310)	S14 321 T-TEST PROBABILITIES	CA, 11-68(115), 13-70(124)
M1	38 A SEARCHING ALGORITHM	CJ, 12-69(101)	S14 322 F-DISTRIBUTION	CA, 11-68(116), 12-69(39)
M1	43 LISTED RADIX SORT	CJ, 12-69(406), 13-70(326)	S14 322	14-71(117)
M1	45 INTERNAL SORT BY TWO-WAY MERGE	CJ, 13-70(110)	S14 344 STUDENT'S T-DISTRIBUTION	CA, 12-69(317), 13-70(124)
M1	55 INTERNAL MERGE SORT	CJ, 13-70(424)	S14 344	13-70(449)
M1	56 TO DISENTANGLE A CHAIN	CJ, 13-70(425)	S14 346 F-TEST PROBABILITIES	CA, 12-69(184)
M1	63 RECURSIVE TREE SORT	CJ, 14-71(103)	S14 349 POLYGAMMA FNS WITH ARB. PRECISION (A)	CA, 12-69(213), 1-75(380)
M1	64 NON-RECURSIVE TREE SORT	CJ, 14-71(104)	S14 349	2-76(206)
M1	INSERTION IN A LIST	JC, 9-62(23)	S14 395 STUDENTS T-DISTRIBUTION	CA, 13-70(617)
M1	SEARCH IN A LIST	JC, 9-62(23)	S14 395 STUDENTS T-DISTRIBUTION (A)	CA, 13-70(617), 5-79(238)
M1	DELETION FROM A LIST	JC, 9-62(24)	S14 396 STUDENTS T-QUANTITIES	CA, 13-70(619)
M1	SORTING WITH MINIMUM STORAGE	JC, 9-62(27)	S14 396 STUDENTS QUANTILES (A)	CA, 13-70(619), 5-79(238)
			S14 404 COMPLEX GAMMA FUNCTION	CA, 14-71(48)
M2	DATA CONVERSION AND SCALING		S14 421 COMPLEX GAMMA FUNCTION	CA, 15-72(271)
M2	DATA PROCESSING-VECTORCARDIOGRM	CA, 5-62(121)	S14 435 MOD. INCOMPLETE GAMMA FUNCTION	CA, 15-72(993)
M2	BINARY MAGNETIC TAPES	CC, 1-70(420)	S14 435 MOD. INCOMPLETE GAMMA FUNCTION (F)	CA, 15-72(993), 4-78(296)
M2	94 HANDLING DATES ON SMALL COMPUTERS (F)	CJ, 20-77(280)	S14 442 NORMAL DEVIAE (A)	CA, 16-73(51)
			S14 465 STUDENT'S T FREQUENCY (A)	CA, 16-73(690)
O2	SIMULATION OF COMPUTING STRUCTURE		S14 487 CUMULATIVE DIST. OF K-S (F)	CA, 17-74(704), 2-76(111)
O2	1 SIM MULTIDIM ARRAY IN ONE DIM	AS, 17-68(180), 18-69(116)	S14 GAMMA AND DIGAMMA FUNCTION (F)	CC, 4-72(221)
O2	18 EVALUATION OF MARGINAL MEANS	AS, 18-69(197)	S14 GAMMA FCN. BY CHEBYSHEV EXP.	NM, 4-63(413)
O2	13 EVALUATION OF FCNAL EXPRESSION	BT, 5-65(137)	S14 518 INCOMPLETE BESSEL FUNC. I SUB 0 (F)	TO, 3-77(279)
O2	100 PROCESSING OF CHAIN-LINKED LIST	CA, 5-62(346)	S14 519 KOLMOGOROV-SMIRNOV PROBABILITIES (F)	TO, 3-77(285)
O2	101 PROCESSING OF CHAIN-LINKED LIST	CA, 5-62(346)	S15 521 RERATED INTEGRALS OF COERROR FUNC. (F)	TO, 3-77(301)
O2	137 NESTED FOR STATEMENT	CA, 5-62(555)	S14 542 INCOMPLETE GAMMA FUNCTION (F)	TO, 5-79()
O2	138 NESTED FOR STATEMENT	CA, 5-62(555)		
O2	SIMULATION OF TURING MACHINE	CP, 8-71(241)		
R2	SYMBOL MANIPULATION		S15 2 NORMAL INTEGRAL	AS, 17-68(186), 18-69(299)
R2	17 BASIC LIST PROCESSING	BT, 6-66(166)	S15 4 AUX FUN FOR DIST INTEGRALS	AS, 17-68(190), 18-69(118)
R2	18 SIMPLIFYING BOOLEAN EXPR (A)	BT, 6-66(260), 12-72(434)	S15 4	19-70(204)
R2	24 FCN. VALUE FROM CHARAC. STRING	BT, 6-69(283)	S15 24 NORMAL INTEGRAL	AS, 18-69(290)
R2	268 ALGOL 60 REF. LANG. EDITOR	CA, 8-65(667), 12-69(407)	S15 66 THE NORMAL INTEGRAL (F)	AS, 22-73(424)
R2	377 SYMB. EXPANSION OF EXPRESSIONS	CA, 13-70(191)	S15 15 COMPL. ERROR INT.-COMPLEX ARG.	BT, 5-65(290)
R2	ALGOL I INTO ALGOL II (A)	CC, 11-76(5)	S15 11 HERMITE POLYNOMIAL	CA, 3-60(353)
R2	CLASSIFICATION OF FORTRAN STMT	CJ, 14-71(100)	S15 123 REAL ERROR FUNCTION	CA, 5-62(483), 6-63(316),
			S15 123	6-63(618), 7-64(145),
			S15 123	10-67(377)
S	APPROXIMATION OF SPECIAL FUNCTIONS...		S15 180 ERROR FUNCTION-LARGE X	CA, 6-63(314), 10-67(377)
S	FUNCTIONS ARE CLASSIFIED S01 TO S22, FOLLOWING		S15 181 COMPLEMENTARY ERR. FCN.-LARGE X	CA, 6-63(315), 7-64(702),
S	FLETCHER-MILLER-ROSENHEAD, INDEX OF MATH. TABLES		S15 181	10-67(377)
S03	19 BINOMIAL COEFFICIENTS	CA, 3-60(540), 5-62(347),	S15 185 ERROR FUNCTION	CA, 6-63(386)
S03	19	5-62(438)	S15 209 ERROR FUNCTION	CA, 6-63(616), 7-64(148),
S03	33 FACTORIAL N	CA, 4-61(106)	S15 226 NORMAL DISTRIBUTION FUNCTION	7-64(482), 10-67(377)
S04	DIRICHLET L-SERIES	MC, 23-69(489)	S15 272 NORMAL DISTRIBUTION FUNCTION	CA, 7-64(295), 10-67(377)
S07	POLAR TRANSF. BY CHEBYSHEV EXP.	NM, 4-63(413)	S15 272	CA, 8-65(789), 10-67(377)
			S15 299 CHI-SQUARED INTEGRAL (A)	11-68(498)
			S15 299	CA, 10-67(243), 11-68(270)
S13	14 COMPLEX EXPONENTIAL INTEGRAL	CA, 3-60(406)	S15 304 NORMAL CURVE INTEGRAL	CA, 10-67(374), 10-67(377)
S13	20 REAL EXPONENTIAL INTEGRAL	CA, 3-60(540), 4-61(105),	S15 304	11-68(271), 12-69(565),
			S15 304	13-70(624)
			S15 363 COMPLEX ERROR FUNCTION	CA, 12-69(635), 15-72(465)
			S15 462 BIVARIATE NORMAL DISTRIB. (F)	CA, 16-73(638)

- S15 DERIV. OF BOYS ERROR FCN. CB, 9-65(105)
S15 13 NORMAL DISTRIBUTION CURVE CJ, 9-66(322), 10-67(113)
S15 39 AREAS UNDER THE NORMAL CURVE CJ, 12-69(197)
S15 ERF(X) BY CHEBYSHEV EXPANSION NM, 4-63(414)
- S16 13 LEGENDRE POLYNOMIAL CA, 3-60(353), 4-61(105),
S16 13 4-61(181)
S16 47 ASSOCIATED LEGENDRE FUNCTION CA, 4-61(178), 6-63(446),
S16 47 12-69(635)
S16 62 ASSOCIATED LEGENDRE FUNCTION CA, 4-61(320), 4-61(544)
S16 259 LEGENDRE FUNCTION CA, 8-65(488)
S16 259 LEGENDRE FUNCS. FOR ARG. .GT. 1 (A) CA, 8-65(488), 3-77(204)
- S17 21 BESSEL FUNCTION CA, 3-60(600), 8-65(219)
S17 22 RICCATI-BESSEL FUNCTION CA, 3-60(600), 13-70(448)
S17 44 BESSEL FUNCTION CA, 4-61(177)
S17 49 SPHERICAL NEUMANN FUNCTION CA, 4-61(179)
S17 49 SPHERICAL NEUMANN FUNCTION (F) CA, 4-61(179), 4-78(295)
S17 124 HANKEL FUNCTION CA, 5-62(483), 8-65(790)
S17 163 HANKEL FUNCTION CA, 6-63(161), 6-63(522)
S17 236 BESSEL FUNCTIONS OF FIRST KIND (A) CA, 7-64(479), 8-65(105),
S17 236 1-75(282)
S17 484 BESSEL K₀, K₁ FOR COMPLEX ARG(F) CA, 17-74(524)
S17 COMPLEX BESSEL FUNCTIONS J₀ AND J₁ (F) CC, 13-77(17)
- S18 3 BESSEL FUNCTION CA, 3-60(240)
S18 0 BESSEL FUNCTION CA, 3-60(240)
S18 214 BESSEL FUNCTION CA, 6-63(662), 7-64(349)
S18 228 Q-BESSEL FUNCTION CA, 7-64(295)
- S18 103 BESSEL FUNCTIONS Y SUB N AND K SUB N (F) CJ, 21-78(272)
S18 511 CDC 6600 BESSEL FUNCTIONS I & J (F) TO, 3-77(93), 4-78()
S18 511 CDC 6600 BESSEL FUNCTIONS I & J (F) TO, 3-77(93)
- S19 57 BERBEI FUNCTION CA, 4-61(181), 5-62(392),
S19 57 5-62(438)
- S20 3 COMPLEMENTARY FRESNEL INTEGRAL BT, 2-62(192)
S20 WEBER FUNCTION BT, 2-62(239)
S20 88 FRESNEL INTEGRALS CA, 5-62(280), 6-63(618)
S20 89 FRESNEL SINE INTEGRAL CA, 5-62(280), 6-63(618)
S20 90 FRESNEL COSINE INTEGRAL CA, 5-62(281), 6-63(618)
S20 213 FRESNEL INTEGRALS CA, 6-63(617), 7-64(661)
S20 244 FRESNEL INTEGRALS CA, 7-64(660)
S20 301 AIRY FUNCTIONS CA, 10-67(291), 10-67(453)
S20 FRESNEL INTEGRALS S(X), C(X) NM, 9-67(382)
S20 505 LIST INSERT. SORT FOR ARB. KEY DIST. (F) TO, 2-76(204)
- S21 53 ELLIPTIC INTEGRAL-FIRST KIND CA, 4-61(180), 6-63(166)
S21 56 ELLIPTIC INTEGRAL-SECOND KIND CA, 4-61(180), 9-66(12)
S21 73 INCOMPLETE ELLIPTIC INTEGRAL CA, 4-61(543), 4-61(544),
S21 73 5-62(514), 6-63(69),
S21 73 6-63(167)
S21 149 ELLIPTIC INTEGRAL CA, 5-62(605), 6-63(166)
S21 149 COMPLETE ELLIPTIC INTEGRAL (A) CA, 5-62(605), 6-63(166),
S21 4-78(95)
S21 165 ELLIPTIC INTEGRAL CA, 6-63(163), 12-69(38)
S21 COMPLETE ELL. INT.-FIRST KIND NM, 5-63(296)
S21 COMPLETE ELL. INT.-SECOND KIND NM, 5-63(297)
S21 COMPLETE ELL. INT. NM, 5-63(297)
S21 INCOMPL. ELL. INT.-FIRST KIND NM, 5-63(297)
S21 INCOMPL. ELL. INT.-SECOND KIND NM, 5-63(298)
S21 INCOMPL. ELL. INT. NM, 5-63(299)
S21 JACOBIAN ELLIPTIC SIN FCN. NM, 5-63(299)
S21 JACOBIAN ELLIPTIC COS FCN. NM, 5-63(300)
S21 JACOBIAN ELLIPTIC FCN. NM, 5-63(301)
S21 ELLIPTIC INTEGRALS-KIND 1, 2, 3 NM, 7-65(85), 7-65(353),
S21 13-69(305)
S21 JACOBIAN ELLIPTIC FUNCTIONS NM, 7-65(89)
S21 COMPLETE ELLIPTIC INTEGRALS (A) NM, 20-73(425)
S21 COMPLETE ELLIPTIC INTEGRALS (A) NM, 20-73(425), 29-78(233)
S21 549 WEIERSTRASS ELLIPTIC FUNCTIONS (F) TO, 6-80(112)
S21 1 ELLIPTIC INTEGRALS ZM, 11-70(99)
- S22 CONFLUENT HYPERG. FCN (COMPLEX) BT, 2-62(237)
S22 5 FERMI FUNCTION BT, 3-63(141)
S22 13 RIEMANN ZETA FUNCTION BT, 5-65(141)
S22 10 CHEBYSHEV POLYNOMIAL CA, 3-60(353), 4-61(181)
S22 12 LAGUERRE POLYNOMIAL CA, 3-60(353)
S22 36 CHEBYSHEV POLYNOMIAL CA, 4-61(151)
S22 110 PHYSICS INTEGRALS CA, 5-62(389), 5-62(393)
S22 111 PHYSICS INTEGRALS CA, 5-62(390)
S22 132 PHYSICS INTEGRALS CA, 5-62(551)
S22 184 ERLANG PROBABILITY FUNCTION CA, 6-63(386)
S22 191 HYPERGEOMETRIC FCN. (COMPLEX) (A) CA, 6-63(388), 7-64(244),
S22 191 17-74(589)
S22 192 CONFLUENT HYPERG. FCN. (COMPLEX) CA, 6-63(388), 7-64(244)
S22 227 CHEBYSHEV POLYNOMIAL COEFF. CA, 7-64(295)
S22 292 REGULAR COULOMB WAVE FCNS. CA, 9-66(793), 12-69(278)
S22 292 12-69(280), 13-70(573)
- S22 300 COULOMB WAVE FUNCTIONS (A) CA, 10-67(244), 12-69(279)
S22 300 12-69(692), 16-73(308)
S22 327 DILOGARITHM CA, 11-68(270)
S22 332 JACOBI POLYNOMIALS (F) CA, 11-68(436), 13-70(449)
S22 332 18-75(116)
S22 352 CHAR. VALS, SLNS OF MATHIEU'S DE. CA, 12-69(399), 13-70(750)
S22 352 15-72(1074)
S22 388 RADEMACHER FCN. CA, 13-70(510)
S22 389 BINARY ORDERED WALSH FCNS. CA, 13-70(511)
S22 390 SEQUENCY ORDERED WALSH FCNS. CA, 13-70(511)
S22 490 DILOGARITHM FUNCTION OF REAL ARGUMENT(F) CA, 18-75(200), 2-76(112)
S22 PHYSICS INTEGRALS CC, 2-71(201)
S22 537 CHAR. VALUES OF MATHIEU'S DIFF. EQN. (F) TO, 5-79(112)
S22 13 CHEBYSHEV POLYNOMIALS ZM, 12-71(227)
S22 282 AND SIN(X)/X 4-69(272), 1-70(53)
- S23 234 POISSON-CHARLIER POLYNOMIALS CA, 7-64(420), 8-65(105)
- Y1 PHYSICS APPLICATIONS
Y1 MANY-ELECTRON WAVE FUNCTIONS CA, 9-66(278)
Y1 INZBURG-LANDAU FLUXOIDS CC, 1-69(10), 1-70(291)
Y1 FRACTIONAL PARENTAGE COEFF CC, 1-69(15)
Y1 FRANK-CONDON FACTORS CC, 1-69(21), 1-70(223)
Y1 COULOMB FUN. FOR COMPLEX ENERGY CC, 1-69(25), 3-72(276)
Y1 RADIATIVE RECOMBINATION COEFF CC, 1-69(31)
Y1 COMPOUND NUCLEAR REACTIONS CC, 1-69(35), 1-70(224)
Y1 NUCLEAR BOUND STATE WAVE FUN CC, 1-69(55)
Y1 MOSSBAUER SPECTRA CC, 1-69(67)
Y1 ELECTRON SCATTERING (F) CC, 1-69(88), 5-73(416)
Y1 REGGE TRAJECTORY CC, 1-69(97)
Y1 NUCLEAR PENETRABILITY CC, 1-69(106)
Y1 CONFIGURATION INTERACTION CC, 1-69(113)
Y1 MULTI-CONF HARTREE-FOCK PROB CC, 1-70(151)
Y1 TWO-ELECTRON WAVE FUNCTIONS CC, 1-70(167)
Y1 HAUSER-FESHBACH NUCLEAR SCAT CC, 1-70(181)
Y1 VECTOR COUPLING COEFFICIENTS CC, 1-70(191)
Y1 NUCLEAR ELASTIC SCATTERING CC, 1-70(198)
Y1 3N-J SYMBOLS FOR SU(2) CC, 1-70(207)
Y1 HARTREE-FOCK-SLATER FIELD CC, 1-70(216)
Y1 ISOSPIN REPRESENTATION CC, 1-70(225)
Y1 COLLISION STRENGTHS CC, 1-70(232)
Y1 RECOUPLING COEFFICIENTS (F) CC, 1-70(241), 2-71(173),
Y1 5-73(161)
Y1 SUBSTITUTION SUMS CC, 1-70(265), 1-70(469)
Y1 MADELUNG POTENTIAL CC, 1-70(280)
Y1 ELECTRON SCATTERING CC, 1-70(306), 1-70(470)
Y1 HYDROGENIC R**K INTEGRALS CC, 1-70(325)
Y1 ANG MOMENTUM COUPLING COEFF CC, 1-70(337)
Y1 SYMMETRIC NEEL WALLS CC, 1-70(342), 1-70(468)
Y1 RAMAN BANDS OF SYMMETRIC TOPS CC, 1-70(349)
Y1 ANGULAR MOMENTUM INTEGRALS CC, 1-70(359)
Y1 CHARGE TRANSFER PROBABILITY CC, 1-70(380)
Y1 GD-STATE WAVE FUN FOR POLYMER CC, 1-70(391)
Y1 MONTE-CARLO GENERATION METHOD CC, 1-70(425)
Y1 INTERACTION MATRIX ELEMENTS CC, 1-70(437)
Y1 SPECTRAL INTENSITY CC, 1-70(440)
Y1 ELECTRON SCATTERING CC, 1-70(445)
Y1 COUPLED ANGULAR MOMENTUM CC, 1-70(453)
Y1 NUMERICAL ORBITAL FUNCTIONS CC, 1-70(457)
Y1 FITTING OF TRANSITION ENERGIES CC, 1-70(465)
Y1 DEPHI-SPEAKEASY SYSTEM CC, 2-71()
Y1 SYMMETRY AND BAND STRUCTURE(1) CC, 2-71(11)
Y1 SYMMETRY AND BAND STRUCTURE(2) CC, 2-71(17)
Y1 SYMMETRY AND BAND STRUCTURE(3) CC, 2-71(26)
Y1 ENERGY LEVEL CALCULATION CC, 2-71(33)
Y1 GAMMA-RAY SPECTRA CC, 2-71(40)
Y1 TRANSPORT COLLISION INTEGRALS CC, 2-71(47)
Y1 ORBIT OF DOUBLE-STAR CC, 2-71(59)
Y1 TRIPLY DEGENERATE RAMAN BANDS CC, 2-71(85)
Y1 BORN APP. FOR NUCLEAR REACTIONS CC, 2-71(94), 3-72(275)
Y1 SELF-CONSISTEN. FIELD (F) CC, 2-71(107), 2-71(471),
Y1 9-75(129)
Y1 CLASSICAL RELATIVE METHOD CC, 2-71(114)
Y1 ELECTRON-ATOM COLLISION CC, 2-71(175)
Y1 ANGULAR MOMENTUM INTEGRALS (F) CC, 2-71(180), 7-74(318),
Y1 8-74(329)
Y1 LANTHANIDE CRYSTAL FIELD CC, 2-71(191)
Y1 PHASE SPACE OF N PARTICLES CC, 2-71(207)
Y1 FOUR-BODY CONTRIBUTIONS CC, 2-71(214)
Y1 PHASE SHIFT AND MIX PARAMETER CC, 2-71(223)
Y1 BRODY-MOSHINSKY BRACKETS CC, 2-71(231)
Y1 ATOMIC WAVEFUNCTION COMP CC, 2-71(239)
Y1 STATIC INTERACTION POTENTIAL(F) CC, 2-71(261), 5-73(396)
Y1 COMPOUND NUCLEAR REACTIONS CC, 2-71(272), 5-73(304)
Y1 GAMMA-RAY DECAY SCHEMES CC, 2-71(288)
Y1 DEGENERATE RAMAN BANDS OF TOPS CC, 2-71(298)
Y1 EXTENDED DEFECTS IN METALS CC, 2-71(301)
Y1 MOSSBAUER SPECTRA CC, 2-71(322)
Y1 PURE AND MIXED NILSSON STATES CC, 2-71(331)
Y1 LEED INTENSITIES CC, 2-71(341)
Y1 TWO-NUCLEON SCHRODINGER EQU. CC, 2-71(353)
Y1 SCATTERING OF ELECTRONS CC, 2-71(360)
Y1 SPECTROSCOPIC DOUBLE STAR CC, 2-71(368)
Y1 ANG. MOMENTUM COUPLING COEFF. CC, 2-71(381)
Y1 BUBBLE CHAMBER PHOTOGRAPH CC, 2-71(394)
Y1 SU3 FRACTIONAL PARENTAGE CC, 2-71(420)
Y1 CONV. COEFF. FOR ATOMIC SHELL CC, 2-71(427)
Y1 ENERGY-LOSS STRAGGLING CC, 2-71(433)

- Y1 NUCLEAR REACTIONS CC, 2-71(443)
- Y1 GAMMA RADIATION DOSIMETRY (F) CC, 2-71(449), 5-73(395), 6-73(240)
- Y1 QUADRUPOLE RADIAL MATRIX CC, 2-71(455)
- Y1 EXTRACTION OF BULK VISCOSITIES CC, 3-72()
- Y1 ATOMIC ENERGY LEVEL VALUES CC, 3-72(11)
- Y1 WAVE NUMBER CALCULATIONS CC, 3-72(19)
- Y1 HARTREE-FOCK NUCL. CALCULATIONS CC, 3-72(22)
- Y1 MAGNETIC SHELL PARAMETER CC, 3-72(31)
- Y1 GEOMAGNETIC SHELL PARAMETER CC, 3-72(36)
- Y1 TRAJECTORY CALCULATIONS CC, 3-72(42)
- Y1 GENZD. TRANSFORMATION BRACKETS CC, 3-72(53)
- Y1 GENERALIZED TALMI COEFFICIENTS CC, 3-72(61)
- Y1 THIRD VIRAL COEFFICIENTS CC, 3-72(69)
- Y1 COULOMB FUNCTIONS CC, 3-72(73)
- Y1 LATTICE VIBRATIONS (F) CC, 3-72(88), 4-72(383)
- Y1 QUADRUPOLE COULOMB EXCITATIONS CC, 3-72(118)
- Y1 FIELD GRADIENT INTEGRALS CC, 3-72(130)
- Y1 BUBBLE CHAMBER EVENTS CC, 3-72(136)
- Y1 CROSS SECTION CALCULATIONS CC, 3-72(173)
- Y1 CLASSICAL RELATIVE MOTION (1) CC, 3-72(197)
- Y1 CLASSICAL RELATIVE MOTION (2) CC, 3-72(221)
- Y1 POSITRON LIFETIME SPECTRA CC, 3-72(240)
- Y1 FINE STRUCTURE CROSS SECTION CC, 3-72(256)
- Y1 TRANSPORT COLLISION INTEGRALS CC, 3-72(269)
- Y1 THE REDUCED ROTATION MATRIX CC, 3-72(318)
- Y1 GROWTH OF GASEOUS DISCHARGE CC, 3-72(322)
- Y1 RADIAL SCHRÖDINGER EQUATION CC, 3-72(334)
- Y1 MOSSBAUER SCATTERING SPECTRA CC, 3-72(339)
- Y1 EINSTEIN TENSOR (F) CC, 4-72(100)
- Y1 MULTI-CONFIG. HARTREE-FOCK (F) CC, 4-72(107), 7-74(236)
- Y1 MULTI-PARTICLE PHASE SPACE (F) CC, 4-72(117)
- Y1 LEVEL-CROSSINGS (F) CC, 4-72(128)
- Y1 LEVEL CROSSINGS & BACK GOUDSMIT EFFECT (F) CC, 4-72(128), 13-77(137)
- Y1 TWO-NUCLEON EFFECTIVE-RANGE (F) CC, 4-72(138)
- Y1 RELATIVISTIC KINEMATICS (F) CC, 4-72(227)
- Y1 RELATIVISTIC KINEMATICS (F) CC, 4-72(233)
- Y1 NUCLEAR WAVE FUNCTION (F) CC, 4-72(239)
- Y1 MOLECULAR VIBRATIONS (F) CC, 4-72(249)
- Y1 NORMALISED MORSE FUNCTIONS (F) CC, 4-72(257)
- Y1 ENERGY LEVEL CALCULATIONS (F) CC, 4-72(262)
- Y1 VECTOR COUPLING COEFFICIENTS (F) CC, 4-72(268)
- Y1 GEOMAGNETIC FIELD MODELS (F) CC, 4-72(347)
- Y1 ENERGY BAND CALCULATIONS (F) CC, 4-72(361)
- Y1 CROSS SECTION & POLARIZATION (F) CC, 4-72(371)
- Y1 CFPJJ-FRACT'L PARENTAGE COEF (F) CC, 4-72(377)
- Y1 CLOSED ORBIT MINIMIZATION (F) CC, 5-73(56)
- Y1 EXPANSION EQUATION OF STATE (F) CC, 5-73(64)
- Y1 NUCLEAR OPTICAL MODEL (F) CC, 5-73(69), 7-74(343)
- Y1 NON-EXCHANGE TYPE INTEGRALS (F) CC, 5-73(80), 8-74(333)
- Y1 EFFECTIVE REGGE TRAJECTORIES (F) CC, 5-73(153)
- Y1 MOMENTUM-SPACE INTEGRALS (A) CC, 5-73(217)
- Y1 MOSSBAUER SPECTRA (F) CC, 5-73(225), 5-73(395)
- Y1 ELECTRON ENERGY DEPOSITION (F) CC, 5-73(239)
- Y1 ANGULAR MOMENTUM COEFFS (F) CC, 5-73(263)
- Y1 HEAVY PARTICLE COLLISIONS (F) CC, 5-73(283)
- Y1 RADIAL DISTRIB. OF EMITTERS (F) CC, 5-73(294)
- Y1 FORM-FACTOR OF LIQ NFE METAL (F) CC, 5-73(299)
- Y1 LIGHT CURVE OF VARIABLE STAR (A) CC, 5-73(315)
- Y1 X-RAY REFLECTION INTENSITIES (F) CC, 5-73(328)
- Y1 ABSORPTION MODEL (F) CC, 5-73(349)
- Y1 SU3 VECTOR COUPLING COEFF (F) CC, 5-73(365)
- Y1 RADIAL SHRODINGER EQUATION (F) CC, 5-73(379)
- Y1 LEGENDRE POLYNOMIALS (F) CC, 5-73(390)
- Y1 SU3 WIGNER & RACAH COEFF (F) CC, 5-73(405)
- Y1 PARTICLE SCATTERING (F) CC, 5-73(417)
- Y1 ELECTROTRANSPORT SIMULATION (F) CC, 5-73(430)
- Y1 AXISYMMETRIC SCALAR FIELD (F) CC, 5-73(437)
- Y1 DIFFERENTIAL CROSS-SECTIONS (F) CC, 5-73(456), 7-74(172)
- Y1 PERTURBATION THEORY (F) CC, 6-73()
- Y1 S-STATE BINDING ENERGY (F) CC, 6-73(17)
- Y1 FREUDENTHAL'S FORMULA (F) CC, 6-73(24)
- Y1 NILSSON ORBITS (F) CC, 6-73(30)
- Y1 RADIATIVE XFER IN SPHER GEOM (F) CC, 6-73(38)
- Y1 ELECTRON-ATOM COLLISION MEAS (F) CC, 6-73(77)
- Y1 A NEW D-SHELL F.P.C. (F) CC, 6-73(88)
- Y1 ATOMIC INTEGRAL CALCULATIONS (F) CC, 6-73(89)
- Y1 COMPOUND NUCLEAR REACTIONS (F) CC, 6-73(99)
- Y1 REDUCED TENSOR-MATRIX ELEMENTS (F) CC, 6-73(132), 9-75(268), 9-75(370)
- Y1 REDUCED TENSOR-MATRIX ELEMENTS (F) CC, 6-73(132), 9-75(268), 9-75(370), 13-77(231)
- Y1 IRREDUCIBLE MULTIPLIER REPRESENTATION (F) CC, 6-73(149), 8-74(141)
- Y1 DYNAM RETRIEV OF ATOM. CODES (F) CC, 6-73(165)
- Y1 ANALYSIS OF REFLECTION DATA (F) CC, 6-73(187)
- Y1 DIATOMIC RKR POTENT'L CURVES (F) CC, 6-73(221)
- Y1 POLAR. EFFECTS IN NUCLR REAC (F) CC, 6-73(229)
- Y1 FINITE-RANGE DWBA CROSS SECT (F) CC, 7-74(13)
- Y1 ELECTRON IMPACT EXCITATION (F) CC, 7-74(38)
- Y1 LEED INTENSITIES CALC. (F) CC, 7-74(50)
- Y1 ENERGIES & SLATER INTEGRALS (F) CC, 7-74(73), 10-75(436)
- Y1 ENERGY DEPOS. BY ION BOMBARD (F) CC, 7-74(85)
- Y1 ENERGY DEPOSITION BY ION BOMB'T (F) CC, 7-74(85), 12-76(335), 12-76(339)
- Y1 RADIATIVE LIFETIMES (F) CC, 7-74(95)
- Y1 ENERGY-LEVEL CALCULATIONS (F) CC, 7-74(145)
- Y1 MOSSBAUER SPECTRA FITTING (F) CC, 7-74(151)
- Y1 CHARGE EXCHANGE PROCESSES (F) CC, 7-74(163)
- Y1 PHOTONS IN 2-LAYERED MEDIA (F) CC, 7-74(185)
- Y1 PHOTONS IN SPHERIC MEDIA (F) CC, 7-74(192)
- Y1 INVER. OF ABEL'S INTEGRAL EQ (F) CC, 7-74(200)
- Y1 CALC. OF CRYSTAL POTENTIALS (F) CC, 7-74(207)
- Y1 LANDAU AND VAVILOV DISTRIB. (F) CC, 7-74(215)
- Y1 MATRIX ELEM OF TENSOR OPTRS (F) CC, 7-74(225)
- Y1 ONE-DIM. LASER FUSION CODE (F) CC, 7-74(271), 10-75(251)
- Y1 THERMAL NEUTRON SCATTERING (F) CC, 7-74(289), 9-75(59)
- Y1 HIGH-ENERGY HADRON CASCADES (F) CC, 7-74(327)
- Y1 SIM. OF EXTENSIVE AIR SHWRS (F) CC, 7-74(344)
- Y1 LEED INTENSITY CURVES (F) CC, 7-74(369)
- Y1 RAYLEIGH FORM FACTORS (F) CC, 7-74(389), 10-75(257)
- Y1 POSITRON LIFETIME SPECTRA (F) CC, 7-74(401)
- Y1 STRUCTURAL DYN. OF POLYTROPIC STARS (F) CC, 7-74(410)
- Y1 WORK FUNCTIONS FROM PHOTOELEC. DATA (F) CC, 7-74(419)
- Y1 COULOMB EXCITATION (F) CC, 8-74(35)
- Y1 SEARCH PROGRAM FOR SIGNIFICANT VARS (F) CC, 8-74(49)
- Y1 GROUP THEORY OF LATTICE DYNAMICS (F) CC, 8-74(71)
- Y1 RACAH'S OUTER MULTIPLICITY FORMULA (F) CC, 8-74(95)
- Y1 WEIZMANN SHELL MODEL (F) CC, 8-74(101)
- Y1 ELASTIC SCATTERING OF PIONS (F) CC, 8-74(130)
- Y1 ATOMIC CONTINUUM PROCESSES (F) CC, 8-74(149)
- Y1 SPIN & PARITY ASSIGNMENTS (F) CC, 8-74(199)
- Y1 BOUND & CONTINUUM ORBITALS (F) CC, 8-74(220)
- Y1 ROTATIONAL BAND SPECTRA (A) CC, 8-74(236)
- Y1 FRACTIONAL PARENTAGE COEFFICIENTS (F) CC, 8-74(246)
- Y1 MULTIPLE "WATER BAG" GRAVITATION SYS (F) CC, 8-74(307)
- Y1 CALCULATIONS FOR NUCLEAR REACTORS (F) CC, 8-74(349)
- Y1 COULOMB WAVE FUNCTIONS (F) CC, 8-74(377), 11-76(141)
- Y1 BOUND STATES OF ONE NUCLEON (F) CC, 8-74(396)
- Y1 COUPLED CHANNEL SCHRÖDINGER EQUATIONS (F) CC, 9-75(11)
- Y1 MULTICONFIGURATION DIRAC-FOCK (F) CC, 9-75(31)
- Y1 MULTICONFIGURATION DERAC-FOCK (F) CC, 9-75(31), 13-77(71)
- Y1 SPECTRA AND CROSS-SECTIONS (F) CC, 9-75(92), 10-75(71)
- Y1 MATRIX OF SPIN-ORBIT INTERACTION (F) CC, 9-75(102), 10-75(70)
- Y1 SPECTRAL CURVE FITTING (F) CC, 9-75(117)
- Y1 CONFIGURE INTERACTION (F) CC, 9-75(141), 10-75(434)
- Y1 ASYMMETRY PARAMETER IN ZEEMAN EFFECT (F) CC, 9-75(173)
- Y1 DIFFRACTIVE EXCITATION MODEL (F) CC, 9-75(182)
- Y1 AXISYMMETRIC VECTOR FIELDS (F) CC, 9-75(193)
- Y1 SPACE GROUP REPRESENTATIONS (F) CC, 9-75(231)
- Y1 SPECT. OF THERMAL RADIOASTRON. SOURCE (F) CC, 9-75(247)
- Y1 RAD. FROM THERM. RADIOASTRON. SOURCES (F) CC, 9-75(257)
- Y1 THE DIRICHLET PROBLEM (F) CC, 9-75(283)
- Y1 CALCS OF CYLINDRICAL PHASE SPACE (F) CC, 9-75(297)
- Y1 COLOUR COORDINATE CALCULATIONS (F) CC, 9-75(305)
- Y1 LEED INTENSITY PATTERNS (F) CC, 9-75(312)
- Y1 CONTINUUM EXCHANGE INTEGRALS (F) CC, 9-75(327)
- Y1 AUTOMATIC SPECTRUM ANALYSIS (F) CC, 9-75(351)
- Y1 REDUCED TENSOR MATRIX ELEMENTS (F) CC, 9-75(370), 13-77(289)
- Y1 OSCILLATOR STRENGTHS--MCHF RADIAL FNS (F) CC, 9-75(381)
- Y1 INT. CONV. COEFFS & PARTICLE PARAMS (F) CC, 9-75(392)
- Y1 EIGEN. OF LABEL. OP. O(3) BASES U(3) (F) CC, 10-75(1)
- Y1 MAGNETOHYDRODYNAMIC STABILITY (F) CC, 10-75(11)
- Y1 X-RAY REFLECTION INTENSITIES (F) CC, 10-75(42)
- Y1 RED. ELEM. SUM. ONE-PARTICLE TENSOR (F) CC, 10-75(56)
- Y1 SYMMETRY AND BANDSTRUCTURE (F) CC, 10-75(67)
- Y1 HARMONIC OSCILLATOR BRACKETS (F) CC, 10-75(87)
- Y1 NORMAL COORDINATE ANAL. OF CRYSTALS (F) CC, 10-75(104)
- Y1 COZ LASER KINETICS CODE (F) CC, 10-75(117)
- Y1 ANAL. OF SPECTROSCOPIC PEAK SHAPES (F) CC, 10-75(145)
- Y1 LASER PULSE PROPAGATION (F) CC, 10-75(155)
- Y1 NUCLEAR POTENTIAL (F) CC, 10-75(182)
- Y1 ATOMIC COLLISIONS (F) CC, 10-75(223)
- Y1 ATOMIC COLLISIONS (F) CC, 10-75(223), 11-76(407)
- Y1 CANTERBURY APPROXIMANTS (A & F) CC, 10-75(234)
- Y1 CATERBURY APPROXIMANTS (A & F) CC, 10-75(234), 13-77(77)
- Y1 CLEBSCH-GORDAN COEFFICIENTS (PF) CC, 10-75(245)
- Y1 DEFORMED QUASIPARTICLES (F) CC, 10-75(293)
- Y1 TRANSFER OF LINE RADIATION (F) CC, 10-75(304)
- Y1 VIBRATIONAL ENERGIES OF CO2 (F) CC, 10-75(368)
- Y1 ROVIBRATIONAL CROSS SECTIONS (F) CC, 10-75(375)
- Y1 KINEMATICS OF THREE-BODY REACTIONS (F) CC, 10-75(385)
- Y1 COULOMB-NUCLEAR INTERFERENCE (F) CC, 10-75(401)
- Y1 MAGNETOTELLURIC MODELLING (F) CC, 10-75(421)
- Y1 ONE-DIM. PULSE HEIGHT SPECTRA (F) CC, 11-76(37)
- Y1 DIPOLE & OVERLAP INTEGRALS (F) CC, 11-76(49)
- Y1 ATOMIC SCF HARTREE-FOCK (A) CC, 11-76(57)
- Y1 GAMMA-GAMMA DIRECTIONAL CORR. COEFF. (F) CC, 11-76(75)
- Y1 INELASTIC SCATTERING OF PIONS (F) CC, 11-76(95)
- Y1 ELEMENTS OF REACTION MATRIX (F) CC, 11-76(113)
- Y1 EVALUATE HYPERFINE STRUCTURE (F) CC, 11-76(125)
- Y1 LAPLACIAN FIELD COMPONENT (F) CC, 11-76(221)
- Y1 STATIC INTERACTION POTENTIAL (F) CC, 11-76(237)
- Y1 QUASI-BOUND STATE WAVEFUNCTIONS (F) CC, 11-76(249)
- Y1 EPR SPIN-HAMILTONIAN PARAMETERS (F) CC, 11-76(257)
- Y1 ANISOTROPIC ELASTIC FIELDS (F) CC, 11-76(279)
- Y1 EVALUATION OF 3J AND 6J COEFFICIENTS (F) CC, 11-76(269)
- Y1 NUCL. REACTS. IN IONOGRAPHIC DETECTORS (F) CC, 11-76(287)
- Y1 CUBIC X-RAY POWDER DEFRACTIONS (F) CC, 11-76(331)
- Y1 ELASTIC GREEN'S TENSOR FUNCTION (A) CC, 11-76(339)
- Y1 ONE-DIM. SOLNS. OF SCHRÖDINGER EQN. (F) CC, 11-76(353)
- Y1 COMPTON SCATTERING OF GAMMA RAYS (F) CC, 11-76(363)
- Y1 TRANSPORT EQN. FOR ELECTRON DISTR. (F) CC, 11-76(369)
- Y1 EQUILIBRIUM POLOIDAL MAGNETIC FIELDS (F) CC, 11-76(385)
- Y1 ANGULAR MOMENTUM COEFFICIENTS (F) CC, 11-76(397)
- Y1 COULOMB COEFFS. FOR IONIC CRYSTALS (A) CC, 12-76(179)
- Y1 ATOMIC COLLISIONS (F) CC, 12-76(199)
- Y1 QUANTUM STATE DENSITIES (F) CC, 12-76(205)
- Y1 1D SIMUL. PLASMA AFTERGLOWS (F) CC, 12-76(213)
- Y1 ENERGY RESPONSE OF SPECTROMETER (F) CC, 12-76(231)
- Y1 MOMENTUM SPACE CODE FOR PIONS (F) CC, 12-76(237)

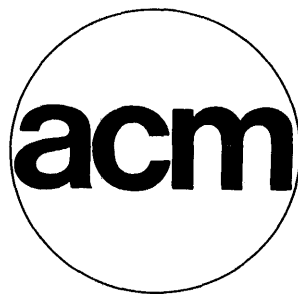
Y1	OPTICAL POTENTIAL CODE FOR PIONS (F)	CC,12-76(237),13-77(141)	Z	391 UNITARY SYMMETRIC POLYNOMIALS	CA,13-70(512),15-72(49)
Y1	SINGLE-PARTICLE-INCLUSIVE PROD. REACT. (F)	CC,12-76(277)	Z	398 TABLELESS DATE CONVERSION	CA,13-70(621),15-72(918)
Y1	TWO-NUCLEON TRANSFER REACTIONS (F)	CC,12-76(293)	Z	428 HU-TUCKER CODING METHOD	CA,15-72(360)
Y1	X-RAY POWDER REFLECTIONS (F)	CC,12-76(305)	Z	444 EXTRACTING PHRASES (A)	CA,16-73(183)
Y1	AMPLIFICATIONS OF LASER LIGHT PULSES (F)	CC,12-76(323)	Z	445 BINARY PATTERN RECONSTRUCT. (A)	CA,16-73(185),16-73(186)
Y1	COUNTING RADIOACTIVE ATOMS (F)	CC,12-77(281)	Z	455 ANAL SKEW REP OF SYMM GROUP (A)	CA,16-73(571)
Y1	CORONAL EMISSION LINE STRENGTHS (F)	CC,13-77(25)	Z	479 MINIMAL SPANNING TREE CLUSTERING METH(F)	CA,17-74(321),18-75(119)
Y1	B SUB N FACTORS FOR EMISSION LINES (F)	CC,13-77(39)	Z	479	2-76(110)
Y1	SHORT RANGE ORDER OF BINARY ALLOYS (F)	CC,13-77(49)	Z	CALCULATION OF EASTER	CB,9-65(18)
Y1	DIFF. CROSS-SECT. FROM TRANS. AMPLI. (F)	CC,13-77(57),13-77(295)	Z	GENERATION OF WEYL GROUP	CC,3-72(155)
Y1	CORIOLIS EFFECT IN ADD-A NUCLEI (F)	CC,13-77(63)	Z	OPIT SYSTEM (M)	CC,5-73(136)
Y1	QUADRUPOLE-DISTORTED NMR POWDER PAT. (F)	CC,13-77(107)	Z	OPTIME SYSTEM (F)	CC,5-73(163)
Y1	LIN. & NONLINEAR IDEAL MHD CODES -V103(P)	CC,13-77(117)	Z	PROGRAM INDEX, VOL. 1-5	CC,5-73(484)
Y1	CORRELATED SELF-DIFFUSION IN CUBIC C. (F)	CC,13-77(167)	Z	OLYMPUS SYSTEM (F)	CC,7-74(237),7-74(245)
Y1	CORRELATION FACTOR AND NMR PAR. IN CC.(F)	CC,13-77(183)	Z		9-75(51),10-75(167)
Y1	ONE BODY SPIN-ORBIT INTERACTION (F)	CC,13-77(193)	Z	MAKING TIMING MEASUREMENTS (F)	CC,8-74(118)
Y1	KRAMERS-KRONIG TRANSFORMS (F)	CC,13-77(203)	Z	OLYMPUS RESTART FACILITIES (F)	CC,8-74(123)
Y1	SIMU. OF POLARIZATION-MOD. ELLIPSO. (F)	CC,13-77(207)	Z	FORMAT AND DATA STATEMENT CONVERTER (F)	CC,11-76(199)
Y1	CALC. OF CRYSTAL POTENTIALS (F)	CC,7-74(207),13-77(225)	Z	FFT PROGRAM GENERATOR (F)	CC,12-76(147)
Y1	COLLISIONS OF PROTONS WITH H ATOMS (F)	CC,13-77(251)	Z	A PORTABLE TEXT EDITOR (F)	CC,13-77(101)
			Z	FORTRAN FUNCTION WRITER (F)	CC,13-77(271)
			Z	SEASONAL ADJ-FORECASTING	CJ,10-67(148),11-68(25)
Z	ALL OTHERS		Z	50 PLAY LEGAL CHESS	CJ,13-70(208),14-71(106)
Z	57 PRINTING MULTIDIM. TABLES (F)	AS,22-73(118)	Z	MULTI-DIMENSION MAP FOLDING	CJ,14-71(25)
Z	25 SYMMETRIC GROUP REPRESENTATION	BT,10-70(106)	Z	68 CHESS GAME	CJ,14-71(209)
Z	DOMINANT WEIGHT IN LIE ALGEBRA	BT,11-71(310)	Z	78 COUNTING VOTES (F)	CJ,16-73(273),14-74(380)
Z	EVALQUOTE IN SIMPLE FORTRAN (F)	BT,12-72(299)	Z	78	18-75(90),18-75(184)
Z	45 INTEREST REFINEMENT	CA,4-61(178),6-63(520)	Z	88 AN ELECTORAL METHOD (A)	CJ,18-75(89),18-75(90)
Z	CALCULATION OF EASTER	CA,5-62(209),5-62(556)	Z	89 REFEREEING KRIEGSPIEL AND CHESS (F)	CJ,18-75(177)
Z	112 POINT INSIDE POLYGON	CA,5-62(434),5-62(606)	Z	BLOCK SIZES FOR MAGNETIC TAPE FILES (A)	CJ,20-77(10)
Z	117 MAGIC SQUARE	CA,5-62(435),5-62(440),	Z	95 GENERATING A PARITY TESTING TABLE (A)	CJ,20-77(280)
Z	117	6-63(39),6-63(105)	Z	105 INTERNAL RATES OF RETURN (F)	CJ,21-78(373),22-79(90)
Z	118 MAGIC SQUARE	CA,5-62(436),5-62(440),	Z	106 UPDATING VISUAL DISPLAY (F)	CJ,22-79(188)
Z	118	5-62(606),6-63(39),	Z	UNIT-CLAUSE PROOF PROCEDURE	CP,7-71(91)
Z	118	6-63(105)	Z	GENERATION OF FINITE GROUPS	CP,7-71(333)
Z	136 ENLARGE A GROUP	CA,5-62(555)	Z	40 INTEGRAL HOMOLOGY FOR TOPOL. GROUPS (P)	CP,23-80(381)
Z	148 MAGIC SQUARE	CA,5-62(605),6-63(168)	Z	COMP. OF GALOIS GROUP ELEMENTS	MC,23-80(425)
Z	199 CALENDAR CONVERSION	CA,6-63(444),7-64(661)	Z	499 AN EFFICIENT SCANNING TECHNIQUE (F)	TO,2-76(82)
Z	240 COORDINATES ON AN ELLIPSOID	CA,7-64(546)	Z	523 CONVEX HULL FOR PLANAR SETS (F)	TO,3-77(411)
Z	246 GRAYCODE (A)	CA,7-64(701),8-65(382),	Z	528 FRAMEWORK FOR A PORTABLE LIBRARY (F)	TO,4-78(177)
Z	246	1-75(285)	Z	528 FRAMEWORK FOR A PORTABLE LIBRARY (F)	TO,4-78(177),5-79()
Z	252 VECTOR COUPLING COEFFICIENTS	CA,8-65(217)	Z	532 SOFTWARE FOR ROUND OFF ANALYSIS (F)	TO,4-78()
Z	GRADER PROGRAM	CA,8-65(277)	Z	536 ONE-WAY ENCRYPTING ALGORITHM (F)	TO,5-79(108)
Z	260 6-J SYMBOLS	CA,8-65(492)	Z	550 SOLID POLYHEDRON MEASURES (F)	TO,6-80(121)
Z	261 9-J SYMBOLS	CA,8-65(492)	Z	561 HEAP PROGRAMS FOR TABLE MAINTENANCE (F)	TO,6-80(444)
Z	355 GENERATE ISING CONFIGURATIONS	CA,12-69(562)	Z	27 PREDICTION OF TIME SERIES (A)	ZM,13-73(553)
Z	364 COLORING POLYGONAL REGIONS	CA,12-69(685)	Z	28 MEASURING ECONOMIC MATURITY (A)	ZM,13-73(559)

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ALGORITHM 493

Zeros of a Real Polynomial [C2]

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Key Words and Phrases: roots, zeros of a polynomial
CR Categories: 5.15
Language: Fortran

DESCRIPTION

The subroutine RPOLY is a Fortran program to find all the zeros of a real polynomial. The parameters are:

OP double precision vector of coefficients in order of decreasing powers of the variable
DEGREE integer degree of the polynomial
ZEROR, double precision vectors of real and imaginary parts of the zeros found
ZEROI by the algorithm
FAIL logical parameter which is true only if the leading coefficient is zero or if RPOLY has found fewer than degree zeros; in the latter case the degree is reset to the number of zeros found.

The routine as written solves polynomials of degree up to 100; however, this can be modified by systematic changing of the declarations in the routines.

The program is based on the three-stage algorithm described in Jenkins and Traub [1]. The algorithm generates a sequence of polynomials of degree one less than the degree of the given polynomial from which an approximation to a zero or a quadratic factor can be extracted. The first stage is linearly convergent and involves no shift of origin. It is used primarily to bias the decision making process in the second stage in favor of the zeros of small magnitude. The second stage is also linearly convergent and involves a double shift to a complex point and its conjugate. The shift point is chosen arbitrarily on a circle whose radius is less than the magnitude of all the zeros. In most cases, either the shift is closest to a real zero, or the pair of shift points are equidistant and closest to a pair of zeros. In the former case the second stage yields an approximation to the real zero and in the latter case it yields an approximation to the real quadratic factor. The third stage involves one of two variable-shift iterations, where the latest approximation(s) is used as the shift point(s) for the generation of the next polynomial. The choice of iteration is based on which case is observed in stage two. The convergence is superquadratic and usually requires only a few steps. The decision processes are made in a fail-safe manner. If the third stage is entered prematurely or the incorrect iteration is chosen, the second stage is resumed. If no convergence is observed in the second stage after a fixed number of steps, a new shift is chosen and the second stage is restarted with a higher fixed limit. The third-stage iterations are terminated when a stopping

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criterion based on roundoff error analysis has been satisfied. The real zero or quadratic factor is removed by polynomial deflation and the algorithm is repeated on the reduced polynomial.

The first statements of RPOLY set the following four constants which describe the floating-point arithmetic of the computer being used:

ETA maximum relative representation error, which can be described as the smallest positive floating-point number such that $1 + \text{ETA} > 1$ in floating-point arithmetic
 INFIN large floating-point number near the top of the range
 SMALNO small positive floating-point number near zero
 BASE exponent base for the floating-point number system.

The program is written in a portable subset of standard Fortran. It has been successfully used on the Burroughs B6700 and the IBM 360/50.

The program has been tested on a large number of polynomials, some chosen to test weaknesses common to zerofinding routines, others randomly generated by a number of techniques.

REFERENCES

1. JENKINS, M.A., AND TRAUB, J.F. A three-stage algorithm for real polynomials using quadratic iteration. *SIAM J. Numer. Anal.* 7 (1970), 545-566.
2. JENKINS, M.A., AND TRAUB, J.F. Principles for testing polynomial zerofinding programs. *ACM TOMS* 1, 1 (March 1975), 26-34.

ALGORITHM

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SUBROUTINE RPOLY(OP, DEGREE, ZEROR, ZEROI,          RPO 10
* FAIL)                                           RPO 20
C FINDS THE ZEROS OF A REAL POLYNOMIAL           RPO 30
C OP - DOUBLE PRECISION VECTOR OF COEFFICIENTS IN RPO 40
C ORDER OF DECREASING POWERS.                   RPO 50
C DEGREE - INTEGER DEGREE OF POLYNOMIAL.         RPO 60
C ZEROR, ZEROI - OUTPUT DOUBLE PRECISION VECTORS OF RPO 70
C REAL AND IMAGINARY PARTS OF THE                RPO 80
C ZEROS.                                          RPO 90
C FAIL - OUTPUT LOGICAL PARAMETER, TRUE ONLY IF RPO 100
C LEADING COEFFICIENT IS ZERO OR IF RPOLY        RPO 110
C HAS FOUND FEWER THAN DEGREE ZEROS.            RPO 120
C IN THE LATTER CASE DEGREE IS RESET TO         RPO 130
C THE NUMBER OF ZEROS FOUND.                    RPO 140
C TO CHANGE THE SIZE OF POLYNOMIALS WHICH CAN BE RPO 150
C SOLVED, RESET THE DIMENSIONS OF THE ARRAYS IN THE RPO 160
C COMMON AREA AND IN THE FOLLOWING DECLARATIONS. RPO 170
C THE SUBROUTINE USES SINGLE PRECISION CALCULATIONS RPO 180
C FOR SCALING, BOUNDS AND ERROR CALCULATIONS. ALL RPO 190
C CALCULATIONS FOR THE ITERATIONS ARE DONE IN DOUBLE RPO 200
C PRECISION.                                     RPO 210
COMMON /GLOBAL/ P, QP, K, QK, SVK, SR, SI, U,    RPO 220
* V, A, B, C, D, A1, A2, A3, A6, A7, E, F, G,    RPO 230
* H, SZR, SZI, LZR, LZI, ETA, ARE, MRE, N, NN    RPO 240
DOUBLE PRECISION P(101), QP(101), K(101),      RPO 250
* QK(101), SVK(101), SR, SI, U, V, A, B, C, D,  RPO 260
* A1, A2, A3, A6, A7, E, F, G, H, SZR, SZI,     RPO 270
* LZR, LZI                                       RPO 280
REAL ETA, ARE, MRE                               RPO 290
INTEGER N, NN                                    RPO 300
DOUBLE PRECISION OP(101), TEMP(101),           RPO 310
* ZEROR(100), ZEROI(100), T, AA, BB, CC, DABS,  RPO 320
* FACTOR                                         RPO 330
REAL PT(101), LO, MAX, MIN, XX, YY, COSR,      RPO 340
* SINR, XXX, X, SC, BND, XM, FF, DF, DX, INFIN, RPO 350
* SMALNO, BASE                                   RPO 360
INTEGER DEGREE, CNT, NZ, I, J, JJ, NMI        RPO 370
LOGICAL FAIL, ZEROK                             RPO 380
C THE FOLLOWING STATEMENTS SET MACHINE CONSTANTS USED RPO 390
C IN VARIOUS PARTS OF THE PROGRAM. THE MEANING OF THE RPO 400
C FOUR CONSTANTS ARE...                          RPO 410
C ETA         THE MAXIMUM RELATIVE REPRESENTATION ERROR RPO 420
C             WHICH CAN BE DESCRIBED AS THE SMALLEST RPO 430
C             POSITIVE FLOATING POINT NUMBER SUCH THAT RPO 440
C              $1.00 + \text{ETA}$  IS GREATER THAN 1. RPO 450
C INFIN       THE LARGEST FLOATING-POINT NUMBER. RPO 460
C SMALNO      THE SMALLEST POSITIVE FLOATING-POINT NUMBER RPO 470
C             IF THE EXPONENT RANGE DIFFERS IN SINGLE AND RPO 480
C             DOUBLE PRECISION THEN SMALNO AND INFIN RPO 490
C             SHOULD INDICATE THE SMALLER RANGE. RPO 500
C BASE        THE BASE OF THE FLOATING-POINT NUMBER RPO 510

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C          SYSTEM USED.                                RPO 520
C THE VALUES BELOW CORRESPOND TO THE BURROUGHS B6700 RPO 530
  BASE = 8.                                           RPO 540
  ETA = .5*BASE**(1-26)                               RPO 550
  INFIN = 4.3E68                                       RPO 560
  SMALNO = 1.0E-45                                     RPO 570
C ARE AND MRE REFER TO THE UNIT ERROR IN + AND *     RPO 580
C RESPECTIVELY. THEY ARE ASSUMED TO BE THE SAME AS  RPO 590
C ETA.                                                RPO 600
  ARE = ETA                                           RPO 610
  MRE = ETA                                           RPO 620
  LO = SMALNO/ETA                                     RPO 630
C INITIALIZATION OF CONSTANTS FOR SHIFT ROTATION     RPO 640
  XX = .70710678                                       RPO 650
  YY = -XX                                           RPO 660
  COSR = -.069756474                                  RPO 670
  SINR = .99756405                                    RPO 680
  FAIL = .FALSE.                                       RPO 690
  N = DEGREE                                          RPO 700
  NN = N + 1                                          RPO 710
C ALGORITHM FAILS IF THE LEADING COEFFICIENT IS ZERO. RPO 720
  IF (OP(1).NE.0.D0) GO TO 10                          RPO 730
  FAIL = .TRUE.                                       RPO 740
  DEGREE = 0                                          RPO 750
  RETURN                                             RPO 760
C REMOVE THE ZEROS AT THE ORIGIN IF ANY              RPO 770
10 IF (OP(NN).NE.0.0D0) GO TO 20                       RPO 780
  J = DEGREE - N + 1                                  RPO 790
  ZEROR(J) = 0.D0                                     RPO 800
  ZEROI(J) = 0.D0                                     RPO 810
  NN = NN - 1                                         RPO 820
  N = N - 1                                           RPO 830
  GO TO 10                                           RPO 840
C MAKE A COPY OF THE COEFFICIENTS                   RPO 850
20 DO 30 I=1,NN                                       RPO 860
  P(I) = OP(I)                                       RPO 870
30 CONTINUE                                          RPO 880
C START THE ALGORITHM FOR ONE ZERO                  RPO 890
40 IF (N.GT.2) GO TO 60                               RPO 900
  IF (N.LT.1) RETURN                                  RPO 910
C CALCULATE THE FINAL ZERO OR PAIR OF ZEROS        RPO 920
  IF (N.EQ.2) GO TO 50                               RPO 930
  ZEROR(DEGREE) = -P(2)/P(1)                         RPO 940
  ZEROI(DEGREE) = 0.0D0                             RPO 950
  RETURN                                             RPO 960
  50 CALL QUAD(P(1), P(2), P(3), ZEROR(DEGREE-1),    RPO 970
  * ZEROI(DEGREE-1), ZEROR(DEGREE), ZEROI(DEGREE))  RPO 980
  RETURN                                             RPO 990
C FIND LARGEST AND SMALLEST MODULI OF COEFFICIENTS. RPO 1000
60 MAX = 0.                                           RPO 1010
  MIN = INFIN                                         RPO 1020
  DO 70 I=1,NN                                       RPO 1030
  X = ABS(SNGL(P(I)))                                 RPO 1040
  IF (X.GT.MAX) MAX = X                               RPO 1050
  IF (X.NE.0. .AND. X.LT.MIN) MIN = X                RPO 1060
70 CONTINUE                                          RPO 1070
C SCALE IF THERE ARE LARGE OR VERY SMALL COEFFICIENTS RPO 1080
C COMPUTES A SCALE FACTOR TO MULTIPLY THE          RPO 1090
C COEFFICIENTS OF THE POLYNOMIAL. THE SCALING IS DONE RPO 1100
C TO AVOID OVERFLOW AND TO AVOID UNDETECTED UNDERFLOW RPO 1110
C INTERFERING WITH THE CONVERGENCE CRITERION.     RPO 1120
C THE FACTOR IS A POWER OF THE BASE                RPO 1130
  SC = LO/MIN                                         RPO 1140
  IF (SC.GT.1.0) GO TO 80                             RPO 1150
  IF (MAX.LT.10.) GO TO 110                           RPO 1160
  IF (SC.EQ.0.) SC = SMALNO                           RPO 1170
  GO TO 90                                           RPO 1180
  80 IF (INFIN/SC.LT.MAX) GO TO 110                   RPO 1190
  90 L = ALOG(SC)/ALOG(BASE) + .5                     RPO 1200
  FACTOR = (BASE*1.0D0)**L                            RPO 1210
  IF (FACTOR.EQ.1.0D0) GO TO 110                     RPO 1220
  DO 100 I=1,NN                                       RPO 1230
  P(I) = FACTOR*P(I)                                  RPO 1240
100 CONTINUE                                          RPO 1250
C COMPUTE LOWER BOUND ON MODULI OF ZEROS.          RPO 1260
110 DO 120 I=1,NN                                     RPO 1270
  PT(I) = ABS(SNGL(P(I)))                             RPO 1280
120 CONTINUE                                          RPO 1290
  PT(NN) = -PT(NN)                                    RPO 1300
C COMPUTE UPPER ESTIMATE OF BOUND                   RPO 1310
  X = EXP((ALOG(-PT(NN))-ALOG(PT(1)))/FLOAT(N))      RPO 1320
  IF (PT(N).EQ.0.) GO TO 130                          RPO 1330
C IF NEWTON STEP AT THE ORIGIN IS BETTER, USE IT.  RPO 1340
  XM = -PT(NN)/PT(N)                                  RPO 1350
  IF (XM.LT.X) X = XM                                 RPO 1360
C CHOP THE INTERVAL (0,X) UNTIL FF .LE. 0         RPO 1370
130 XM = X*.1                                         RPO 1380
  FF = PT(1)                                          RPO 1390
  DO 140 I=2,NN                                       RPO 1400
  FF = FF*XM + PT(I)                                  RPO 1410
140 CONTINUE                                          RPO 1420
  IF (FF.LE.0.) GO TO 150                            RPO 1430

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X = XM
GO TO 130
150 DX = X
C DO NEWTON ITERATION UNTIL X CONVERGES TO TWO
C DECIMAL PLACES
160 IF (ABS(DX/X).LE..005) GO TO 180
FF = PT(1)
DF = FF
DO 170 I=2,N
FF = FF*X + PT(I)
DF = DF*X + FF
170 CONTINUE
FF = FF*X + PT(NN)
DX = FF/DF
X = X - DX
GO TO 160
180 BND = X
C COMPUTE THE DERIVATIVE AS THE INITIAL K POLYNOMIAL
C AND DO 5 STEPS WITH NO SHIFT
NM1 = N - 1
DO 190 I=2,N
K(I) = FLOAT(NN-I)*P(I)/FLOAT(N)
190 CONTINUE
K(1) = P(1)
AA = P(NN)
BB = P(N)
ZEROK = K(N).EQ.0.D0
DO 230 JJ=1,5
CC = K(N)
IF (ZEROK) GO TO 210
C USE SCALED FORM OF RECURRENCE IF VALUE OF K AT 0 IS
C NONZERO
T = -AA/CC
DO 200 I=1,NM1
J = NN - I
K(J) = T*K(J-1) + P(J)
200 CONTINUE
K(1) = P(1)
ZEROK = DABS(K(N)).LE.DABS(BB)*ETA*10.
GO TO 230
C USE UNSCALED FORM OF RECURRENCE
210 DO 220 I=1,NM1
J = NN - I
K(J) = K(J-1)
220 CONTINUE
K(1) = 0.D0
ZEROK = K(N).EQ.0.D0
230 CONTINUE
C SAVE K FOR RESTARTS WITH NEW SHIFTS
DO 240 I=1,N
TEMP(I) = K(I)
240 CONTINUE
C LOOP TO SELECT THE QUADRATIC CORRESPONDING TO EACH
C NEW SHIFT
DO 280 CNT=1,20
C QUADRATIC CORRESPONDS TO A DOUBLE SHIFT TO A
C NON-REAL POINT AND ITS COMPLEX CONJUGATE. THE POINT
C HAS MODULUS BND AND AMPLITUDE ROTATED BY 94 DEGREES
C FROM THE PREVIOUS SHIFT
XXX = COSR*XX - SINR*YY
YY = SINR*XX + COSR*YY
XX = XXX
SR = BND*XX
SI = BND*YY
U = -2.0D0*SR
V = BND
C SECOND STAGE CALCULATION, FIXED QUADRATIC
CALL FXSHFR(20*CNT, NZ)
IF (NZ.EQ.0) GO TO 260
C THE SECOND STAGE JUMPS DIRECTLY TO ONE OF THE THIRD
C STAGE ITERATIONS AND RETURNS HERE IF SUCCESSFUL.
C DEFLATE THE POLYNOMIAL, STORE THE ZERO OR ZEROS AND
C RETURN TO THE MAIN ALGORITHM.
J = DEGREE - N + 1
ZEROR(J) = SZR
ZEROI(J) = SZI
NN = NN - NZ
N = NN - 1
DO 250 I=1,NN
P(I) = QP(I)
250 CONTINUE
IF (NZ.EQ.1) GO TO 40
ZEROR(J+1) = LZR
ZEROI(J+1) = LZI
GO TO 40
C IF THE ITERATION IS UNSUCCESSFUL ANOTHER QUADRATIC
C IS CHOSEN AFTER RESTORING K
260 DO 270 I=1,N
K(I) = TEMP(I)
270 CONTINUE
280 CONTINUE
C RETURN WITH FAILURE IF NO CONVERGENCE WITH 20

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RPO 1440
RPO 1450
RPO 1460
RPO 1470
RPO 1480
RPO 1490
RPO 1500
RPO 1510
RPO 1520
RPO 1530
RPO 1540
RPO 1550
RPO 1560
RPO 1570
RPO 1580
RPO 1590
RPO 1600
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RPO 2110
RPO 2120
RPO 2130
RPO 2140
RPO 2150
RPO 2160
RPO 2170
RPO 2180
RPO 2190
RPO 2200
RPO 2210
RPO 2220
RPO 2230
RPO 2240
RPO 2250
RPO 2260
RPO 2270
RPO 2280
RPO 2290
RPO 2300
RPO 2310
RPO 2320
RPO 2330
RPO 2340
RPO 2350

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C SHIFTS	RPO 2360
FAIL = .TRUE.	RPO 2370
DEGREE = DEGREE - N	RPO 2380
RETURN	RPO 2390
END	RPO 2400
SUBROUTINE FXSHFR(L2, NZ)	FXS 10
C COMPUTES UP TO L2 FIXED SHIFT K-POLYNOMIALS,	FXS 20
C TESTING FOR CONVERGENCE IN THE LINEAR OR QUADRATIC	FXS 30
C CASE. INITIATES ONE OF THE VARIABLE SHIFT	FXS 40
C ITERATIONS AND RETURNS WITH THE NUMBER OF ZEROS	FXS 50
C FOUND.	FXS 60
C L2 - LIMIT OF FIXED SHIFT STEPS	FXS 70
C NZ - NUMBER OF ZEROS FOUND	FXS 80
COMMON /GLOBAL/ P, QP, K, QK, SVK, SR, SI, U,	FXS 90
* V, A, B, C, D, A1, A2, A3, A6, A7, E, F, G,	FXS 100
* H, SZR, SZI, LZR, LZI, ETA, ARE, MRE, N, NN	FXS 110
DOUBLE PRECISION P(101), QP(101), K(101),	FXS 120
* QK(101), SVK(101), SR, SI, U, V, A, B, C, D,	FXS 130
* A1, A2, A3, A6, A7, E, F, G, H, SZR, SZI,	FXS 140
* LZR, LZI	FXS 150
REAL ETA, ARE, MRE	FXS 160
INTEGER N, NN	FXS 170
DOUBLE PRECISION SVU, SVV, UI, VI, S	FXS 180
REAL BETAS, BETAV, OSS, OVV, SS, VV, TS, TV,	FXS 190
* OTS, OTV, TVV, TSS	FXS 200
INTEGER L2, NZ, TYPE, I, J, IFLAG	FXS 210
LOGICAL VPASS, SPASS, VTRY, STRY	FXS 220
NZ = 0	FXS 230
BETAV = .25	FXS 240
BETAS = .25	FXS 250
OSS = SR	FXS 260
OVV = V	FXS 270
C EVALUATE POLYNOMIAL BY SYNTHETIC DIVISION	FXS 280
CALL QUADSD(NN, U, V, P, QP, A, B)	FXS 290
CALL CALCSC(TYPE)	FXS 300
DO 80 J=1,L2	FXS 310
C CALCULATE NEXT K POLYNOMIAL AND ESTIMATE V	FXS 320
CALL NEXTK(TYPE)	FXS 330
CALL CALCSC(TYPE)	FXS 340
CALL NEWEST(TYPE, UI, VI)	FXS 350
VV = VI	FXS 360
C ESTIMATE S	FXS 370
SS = 0.	FXS 380
IF (K(N).NE.0.D0) SS = -P(NN)/K(N)	FXS 390
TV = 1.	FXS 400
TS = 1.	FXS 410
IF (J.EQ.1 .OR. TYPE.EQ.3) GO TO 70	FXS 420
C COMPUTE RELATIVE MEASURES OF CONVERGENCE OF S AND V	FXS 430
C SEQUENCES	FXS 440
IF (VV.NE.0.) TV = ABS((VV-OVV)/VV)	FXS 450
IF (SS.NE.0.) TS = ABS((SS-OSS)/SS)	FXS 460
C IF DECREASING, MULTIPLY TWO MOST RECENT	FXS 470
C CONVERGENCE MEASURES	FXS 480
TVV = 1.	FXS 490
IF (TV.LT.OTV) TVV = TV*OTV	FXS 500
TSS = 1.	FXS 510
IF (TS.LT.OTS) TSS = TS*OTS	FXS 520
C COMPARE WITH CONVERGENCE CRITERIA	FXS 530
VPASS = TVV.LT.BETAV	FXS 540
SPASS = TSS.LT.BETAS	FXS 550
IF (.NOT.(SPASS .OR. VPASS)) GO TO 70	FXS 560
C AT LEAST ONE SEQUENCE HAS PASSED THE CONVERGENCE	FXS 570
C TEST. STORE VARIABLES BEFORE ITERATING	FXS 580
SVU = U	FXS 590
SVV = V	FXS 600
DO 10 I=1,N	FXS 610
SVK(I) = K(I)	FXS 620
10 CONTINUE	FXS 630
S = SS	FXS 640
C CHOOSE ITERATION ACCORDING TO THE FASTEST	FXS 650
C CONVERGING SEQUENCE	FXS 660
VTRY = .FALSE.	FXS 670
STRY = .FALSE.	FXS 680
IF (SPASS .AND. ((.NOT.VPASS) .OR.	FXS 690
* TSS.LT.TVV)) GO TO 40	FXS 700
20 CALL QUADIT(UI, VI, NZ)	FXS 710
IF (NZ.GT.0) RETURN	FXS 720
C QUADRATIC ITERATION HAS FAILED. FLAG THAT IT HAS	FXS 730
C BEEN TRIED AND DECREASE THE CONVERGENCE CRITERION.	FXS 740
VTRY = .TRUE.	FXS 750
BETAV = BETAV*.25	FXS 760
C TRY LINEAR ITERATION IF IT HAS NOT BEEN TRIED AND	FXS 770
C THE S SEQUENCE IS CONVERGING	FXS 780
IF (STRY .OR. (.NOT.SPASS)) GO TO 50	FXS 790
DO 30 I=1,N	FXS 800
K(I) = SVK(I)	FXS 810
30 CONTINUE	FXS 820
40 CALL REALIT(S, NZ, IFLAG)	FXS 830
IF (NZ.GT.0) RETURN	FXS 840

```

C LINEAR ITERATION HAS FAILED. FLAG THAT IT HAS BEEN          FXS 850
C TRIED AND DECREASE THE CONVERGENCE CRITERION              FXS 860
    STRY = .TRUE.                                           FXS 870
    BETAS = BETAS*.25                                       FXS 880
    IF (IFLAG.EQ.0) GO TO 50                                FXS 890
C IF LINEAR ITERATION SIGNALS AN ALMOST DOUBLE REAL        FXS 900
C ZERO ATTEMPT QUADRATIC INTERATION                         FXS 910
    UI = -(S+S)                                             FXS 920
    VI = S*S                                                FXS 930
    GO TO 20                                                FXS 940
C RESTORE VARIABLES                                         FXS 950
50  U = SVU                                                 FXS 960
    V = SVV                                                 FXS 970
    DO 60 I=1,N                                             FXS 980
        K(I) = SVK(I)                                       FXS 990
    60  CONTINUE                                           FXS 1000
C TRY QUADRATIC ITERATION IF IT HAS NOT BEEN TRIED        FXS 1010
C AND THE V SEQUENCE IS CONVERGING                         FXS 1020
    IF (VPASS .AND. (.NOT.VTRY)) GO TO 20                 FXS 1030
C RECOMPUTE QP AND SCALAR VALUES TO CONTINUE THE         FXS 1040
C SECOND STAGE                                             FXS 1050
    CALL QUADSD(NN, U, V, P, QP, A, B)                    FXS 1060
    CALL CALCSC(TYPE)                                     FXS 1070
70  OVV = VV                                               FXS 1080
    OSS = SS                                               FXS 1090
    OTV = TV                                               FXS 1100
    OTS = TS                                               FXS 1110
80  CONTINUE                                               FXS 1120
    RETURN                                                 FXS 1130
    END                                                    FXS 1140

    SUBROUTINE QUADIT(UU, VV, NZ)
C VARIABLE-SHIFT K-POLYNOMIAL ITERATION FOR A
C QUADRATIC FACTOR CONVERGES ONLY IF THE ZEROS ARE
C EQUIMODULAR OR NEARLY SO.
C UU,VV - COEFFICIENTS OF STARTING QUADRATIC
C NZ - NUMBER OF ZERO FOUND
    COMMON /GLOBAL/ P, QP, K, QK, SVK, SR, SI, U,
    * V, A, B, C, D, A1, A2, A3, A6, A7, E, F, G,
    * H, SZR, SZI, LZR, LZI, ETA, ARE, MRE, N, NN
    DOUBLE PRECISION P(101), QP(101), K(101),
    * QK(101), SVK(101), SR, SI, U, V, A, B, C, D,
    * A1, A2, A3, A6, A7, E, F, G, H, SZR, SZI,
    * LZR, LZI
    REAL ETA, ARE, MRE
    INTEGER N, NN
    DOUBLE PRECISION UI, VI, UU, VV, DABS
    REAL MS, MP, OMP, EE, RELSTP, T, ZM
    INTEGER NZ, TYPE, I, J
    LOGICAL TRIED
    NZ = 0
    TRIED = .FALSE.
    U = UU
    V = VV
    J = 0
C MAIN LOOP
10  CALL QUAD(1.D0, U, V, SZR, SZI, LZR, LZI)
C RETURN IF ROOTS OF THE QUADRATIC ARE REAL AND NOT
C CLOSE TO MULTIPLE OR NEARLY EQUAL AND OF OPPOSITE
C SIGN
    IF (DABS(DABS(SZR)-DABS(LZR)).GT..01D0*
    * DABS(LZR)) RETURN
C EVALUATE POLYNOMIAL BY QUADRATIC SYNTHETIC DIVISION
    CALL QUADSD(NN, U, V, P, QP, A, B)
    MP = DABS(A-SZR*B) + DABS(SZI*B)
C COMPUTE A RIGOROUS BOUND ON THE ROUNDING ERROR IN
C EVALUTING P
    ZM = SQRT(ABS(SNGL(V)))
    EE = 2.*ABS(SNGL(QP(1)))
    T = -SZR*B
    DO 20 I=2,N
        EE = EE*ZM + ABS(SNGL(QP(I)))
    20  CONTINUE
    EE = EE*ZM + ABS(SNGL(A)+T)
    EE = (5.*MRE+4.*ARE)*EE - (5.*MRE+2.*ARE)*
    * (ABS(SNGL(A)+T)+ABS(SNGL(B))*ZM) +
    * 2.*ARE*ABS(T)
C ITERATION HAS CONVERGED SUFFICIENTLY IF THE
C POLYNOMIAL VALUE IS LESS THAN 20 TIMES THIS BOUND
    IF (MP.GT.20.*EE) GO TO 30
    NZ = 2
    RETURN
    30  J = J + 1
C STOP ITERATION AFTER 20 STEPS
    IF (J.GT.20) RETURN
    IF (J.LT.2) GO TO 50
    IF (RELSTP.GT..01 .OR. MP.LT.OMP .OR. TRIED)
    * GO TO 50
C A CLUSTER APPEARS TO BE STALLING THE CONVERGENCE.

```

```

QUA 10
QUA 20
QUA 30
QUA 40
QUA 50
QUA 60
QUA 70
QUA 80
QUA 90
QUA 100
QUA 110
QUA 120
QUA 130
QUA 140
QUA 150
QUA 160
QUA 170
QUA 180
QUA 190
QUA 200
QUA 210
QUA 220
QUA 230
QUA 240
QUA 250
QUA 260
QUA 270
QUA 280
QUA 290
QUA 300
QUA 310
QUA 320
QUA 330
QUA 340
QUA 350
QUA 360
QUA 370
QUA 380
QUA 390
QUA 400
QUA 410
QUA 420
QUA 430
QUA 440
QUA 450
QUA 460
QUA 470
QUA 480
QUA 490
QUA 500
QUA 510
QUA 520
QUA 530
QUA 540
QUA 550
QUA 560
QUA 570
QUA 580

```

C FIVE FIXED SHIFT STEPS ARE TAKEN WITH A U,V CLOSE	QUA 590
C TO THE CLUSTER	QUA 600
IF (RELSTP.LT.ETA) RELSTP = ETA	QUA 610
RELSTP = SQRT(RELSTP)	QUA 620
U = U - U*RELSTP	QUA 630
V = V + V*RELSTP	QUA 640
CALL QUADSD(NN, U, V, P, QP, A, B)	QUA 650
DO 40 I=1,5	QUA 660
CALL CALCSC(TYPE)	QUA 670
CALL NEXTK(TYPE)	QUA 680
40 CONTINUE	QUA 690
TRIED = .TRUE.	QUA 700
J = 0	QUA 710
50 OMP = MP	QUA 720
C CALCULATE NEXT K POLYNOMIAL AND NEW U AND V	QUA 730
CALL CALCSC(TYPE)	QUA 740
CALL NEXTK(TYPE)	QUA 750
CALL CALCSC(TYPE)	QUA 760
CALL NEWEST(TYPE, UI, VI)	QUA 770
C IF VI IS ZERO THE ITERATION IS NOT CONVERGING	QUA 780
IF (VI.EQ.0.D0) RETURN	QUA 790
RELSTP = DABS((VI-V)/VI)	QUA 800
U = UI	QUA 810
V = VI	QUA 820
GO TO 10	QUA 830
END	QUA 840
SUBROUTINE REALIT(SSS, NZ, IFLAG)	REA 10
C VARIABLE-SHIFT H POLYNOMIAL ITERATION FOR A REAL	REA 20
C ZERO.	REA 30
C SSS - STARTING ITERATE	REA 40
C NZ - NUMBER OF ZERO FOUND	REA 50
C IFLAG - FLAG TO INDICATE A PAIR OF ZEROS NEAR REAL	REA 60
C AXIS.	REA 70
COMMON /GLOBAL/ P, QP, K, QK, SVK, SR, SI, U,	REA 80
* V, A, B, C, D, A1, A2, A3, A6, A7, E, F, G,	REA 90
* H, SZR, SZI, LZR, LZI, ETA, ARE, MRE, N, NN	REA 100
DOUBLE PRECISION P(101), QP(101), K(101),	REA 110
* QK(101), SVK(101), SR, SI, U, V, A, B, C, D,	REA 120
* A1, A2, A3, A6, A7, E, F, G, H, SZR, SZI,	REA 130
* LZR, LZI	REA 140
REAL ETA, ARE, MRE	REA 150
INTEGER N, NN	REA 160
DOUBLE PRECISION PV, KV, T, S, SSS, DABS	REA 170
REAL MS, MP, OMP, EE	REA 180
INTEGER NZ, IFLAG, I, J, NMI	REA 190
NMI = N - 1	REA 200
NZ = 0	REA 210
S = SSS	REA 220
IFLAG = 0	REA 230
J = 0	REA 240
C MAIN LOOP	REA 250
10 PV = P(1)	REA 260
C EVALUATE P AT S	REA 270
QP(1) = PV	REA 280
DO 20 I=2,NN	REA 290
PV = PV*S + P(I)	REA 300
QP(I) = PV	REA 310
20 CONTINUE	REA 320
MP = DABS(PV)	REA 330
C COMPUTE A RIGOROUS BOUND ON THE ERROR IN EVALUATING	REA 340
C P	REA 350
MS = DABS(S)	REA 360
EE = (MRE/(ARE+MRE))*ABS(SNGL(QP(1)))	REA 370
DO 30 I=2,NN	REA 380
EE = EE*MS + ABS(SNGL(QP(I)))	REA 390
30 CONTINUE	REA 400
C ITERATION HAS CONVERGED SUFFICIENTLY IF THE	REA 410
C POLYNOMIAL VALUE IS LESS THAN 20 TIMES THIS BOUND	REA 420
IF (MP.GT.20.*((ARE+MRE)*EE-MRE*MP)) GO TO 40	REA 430
NZ = 1	REA 440
SZR = S	REA 450
SZI = 0.D0	REA 460
RETURN	REA 470
40 J = J + 1	REA 480
C STOP ITERATION AFTER 10 STEPS	REA 490
IF (J.GT.10) RETURN	REA 500
IF (J.LT.2) GO TO 50	REA 510
IF (DABS(T).GT..001*DABS(S-T) .OR. MP.LE.OMP)	REA 520
* GO TO 50	REA 530
C A CLUSTER OF ZEROS NEAR THE REAL AXIS HAS BEEN	REA 540
C ENCOUNTERED RETURN WITH IFLAG SET TO INITIATE A	REA 550
C QUADRATIC ITERATION	REA 560
IFLAG = 1	REA 570
SSS = S	REA 580
RETURN	REA 590
C RETURN IF THE POLYNOMIAL VALUE HAS INCREASED	REA 600
C SIGNIFICANTLY	REA 610
50 OMP = MP	REA 620

```

C COMPUTE T, THE NEXT POLYNOMIAL, AND THE NEW ITERATE
  KV = K(1)
  QK(1) = KV
  DO 60 I=2,N
    KV = KV*S + K(I)
    QK(I) = KV
  60 CONTINUE
  IF (DABS(KV).LE.DABS(K(N))*10.*ETA) GO TO 80
C USE THE SCALED FORM OF THE RECURRENCE IF THE VALUE
C OF K AT S IS NONZERO
  T = -PV/KV
  K(1) = QP(1)
  DO 70 I=2,N
    K(I) = T*QK(I-1) + QP(I)
  70 CONTINUE
  GO TO 100
C USE UNSCALED FORM
  80 K(1) = 0.0D0
  DO 90 I=2,N
    K(I) = QK(I-1)
  90 CONTINUE
  100 KV = K(1)
  DO 110 I=2,N
    KV = KV*S + K(I)
  110 CONTINUE
  T = 0.D0
  IF (DABS(KV).GT.DABS(K(N))*10.*ETA) T = -PV/KV
  S = S + T
  GO TO 10
END

```

REA 630
 REA 640
 REA 650
 REA 660
 REA 670
 REA 680
 REA 690
 REA 700
 REA 710
 REA 720
 REA 730
 REA 740
 REA 750
 REA 760
 REA 770
 REA 780
 REA 790
 REA 800
 REA 810
 REA 820
 REA 830
 REA 840
 REA 850
 REA 860
 REA 870
 REA 880
 REA 890
 REA 900
 REA 910
 REA 920

```

SUBROUTINE CALCSC(TYPE)
C THIS ROUTINE CALCULATES SCALAR QUANTITIES USED TO
C COMPUTE THE NEXT K POLYNOMIAL AND NEW ESTIMATES OF
C THE QUADRATIC COEFFICIENTS.
C TYPE - INTEGER VARIABLE SET HERE INDICATING HOW THE
C CALCULATIONS ARE NORMALIZED TO AVOID OVERFLOW
COMMON /GLOBAL/ P, QP, K, QK, SVK, SR, SI, U,
* V, A, B, C, D, A1, A2, A3, A6, A7, E, F, G,
* H, SZR, SZI, LZR, LZI, ETA, ARE, MRE, N, NN
DOUBLE PRECISION P(101), QP(101), K(101),
* QK(101), SVK(101), SR, SI, U, V, A, B, C, D,
* A1, A2, A3, A6, A7, E, F, G, H, SZR, SZI,
* LZR, LZI
REAL ETA, ARE, MRE
INTEGER N, NN
DOUBLE PRECISION DABS
INTEGER TYPE
C SYNTHETIC DIVISION OF K BY THE QUADRATIC 1,U,V
CALL QUADSD(N, U, V, K, QK, C, D)
IF (DABS(C).GT.DABS(K(N))*100.*ETA) GO TO 10
IF (DABS(D).GT.DABS(K(N-1))*100.*ETA) GO TO 10
TYPE = 3
C TYPE=3 INDICATES THE QUADRATIC IS ALMOST A FACTOR
C OF K
RETURN
10 IF (DABS(D).LT.DABS(C)) GO TO 20
TYPE = 2
C TYPE=2 INDICATES THAT ALL FORMULAS ARE DIVIDED BY D
E = A/D
F = C/D
G = U*B
H = V*B
A3 = (A+G)*E + H*(B/D)
A1 = B*F - A
A7 = (F+U)*A + H
RETURN
20 TYPE = 1
C TYPE=1 INDICATES THAT ALL FORMULAS ARE DIVIDED BY C
E = A/C
F = D/C
G = U*E
H = V*B
A3 = A*E + (H/C+G)*B
A1 = B - A*(D/C)
A7 = A + G*D + H*F
RETURN
END

```

CAL 10
 CAL 20
 CAL 30
 CAL 40
 CAL 50
 CAL 60
 CAL 70
 CAL 80
 CAL 90
 CAL 100
 CAL 110
 CAL 120
 CAL 130
 CAL 140
 CAL 150
 CAL 160
 CAL 170
 CAL 180
 CAL 190
 CAL 200
 CAL 210
 CAL 220
 CAL 230
 CAL 240
 CAL 250
 CAL 260
 CAL 270
 CAL 280
 CAL 290
 CAL 300
 CAL 310
 CAL 320
 CAL 330
 CAL 340
 CAL 350
 CAL 360
 CAL 370
 CAL 380
 CAL 390
 CAL 400
 CAL 410
 CAL 420
 CAL 430
 CAL 440
 CAL 450
 CAL 460
 CAL 470

```

SUBROUTINE NEXTK(TYPE)
C COMPUTES THE NEXT K POLYNOMIALS USING SCALARS
C COMPUTED IN CALCSC
COMMON /GLOBAL/ P, QP, K, QK, SVK, SR, SI, U,
* V, A, B, C, D, A1, A2, A3, A6, A7, E, F, G,
* H, SZR, SZI, LZR, LZI, ETA, ARE, MRE, N, NN
DOUBLE PRECISION P(101), QP(101), K(101),

```

NEX 10
 NEX 20
 NEX 30
 NEX 40
 NEX 50
 NEX 60
 NEX 70

```

* QK(101), SVK(101), SR, SI, U, V, A, B, C, D,
* A1, A2, A3, A6, A7, E, F, G, H, SZR, SZI,
* LZR, LZI
REAL ETA, ARE, MRE
INTEGER N, NN
DOUBLE PRECISION TEMP, DABS
INTEGER TYPE
IF (TYPE.EQ.3) GO TO 40
TEMP = A
IF (TYPE.EQ.1) TEMP = B
IF (DABS(A1).GT.DABS(TEMP)*ETA*10.) GO TO 20
C IF A1 IS NEARLY ZERO THEN USE A SPECIAL FORM OF THE
C RECURRENCE
K(1) = 0.D0
K(2) = -A7*QP(1)
DO 10 I=3,N
K(I) = A3*QK(I-2) - A7*QP(I-1)
10 CONTINUE
RETURN
C USE SCALED FORM OF THE RECURRENCE
20 A7 = A7/A1
A3 = A3/A1
K(1) = QP(1)
K(2) = QP(2) - A7*QP(1)
DO 30 I=3,N
K(I) = A3*QK(I-2) - A7*QP(I-1) + QP(I)
30 CONTINUE
RETURN
C USE UNSCALED FORM OF THE RECURRENCE IF TYPE IS 3
40 K(1) = 0.D0
K(2) = 0.D0
DO 50 I=3,N
K(I) = QK(I-2)
50 CONTINUE
RETURN
END
NEX 80
NEX 90
NEX 100
NEX 110
NEX 120
NEX 130
NEX 140
NEX 150
NEX 160
NEX 170
NEX 180
NEX 190
NEX 200
NEX 210
NEX 220
NEX 230
NEX 240
NEX 250
NEX 260
NEX 270
NEX 280
NEX 290
NEX 300
NEX 310
NEX 320
NEX 330
NEX 340
NEX 350
NEX 360
NEX 370
NEX 380
NEX 390
NEX 400
NEX 410
NEX 420
NEX 430

SUBROUTINE NEWEST(TYPE, UU, VV)
C COMPUTE NEW ESTIMATES OF THE QUADRATIC COEFFICIENTS
C USING THE SCALARS COMPUTED IN CALCSC.
COMMON /GLOBAL/ P, QP, K, QK, SVK, SR, SI, U,
* V, A, B, C, D, A1, A2, A3, A6, A7, E, F, G,
* H, SZR, SZI, LZR, LZI, ETA, ARE, MRE, N, NN
DOUBLE PRECISION P(101), QP(101), K(101),
* QK(101), SVK(101), SR, SI, U, V, A, B, C, D,
* A1, A2, A3, A6, A7, E, F, G, H, SZR, SZI,
* LZR, LZI
REAL ETA, ARE, MRE
INTEGER N, NN
DOUBLE PRECISION A4, A5, B1, B2, C1, C2, C3,
* C4, TEMP, UU, VV
INTEGER TYPE
C USE FORMULAS APPROPRIATE TO SETTING OF TYPE.
IF (TYPE.EQ.3) GO TO 30
IF (TYPE.EQ.2) GO TO 10
A4 = A + U*B + H*F
A5 = C + (U+V*F)*D
GO TO 20
10 A4 = (A+G)*F + H
A5 = (F+U)*C + V*D
C EVALUATE NEW QUADRATIC COEFFICIENTS.
20 B1 = -K(N)/P(NN)
B2 = -(K(N-1)+B1*P(N))/P(NN)
C1 = V*B2*A1
C2 = B1*A7
C3 = B1*B1*A3
C4 = C1 - C2 - C3
TEMP = A5 + B1*A4 - C4
IF (TEMP.EQ.0.D0) GO TO 30
UU = U - (U*(C3+C2)+V*(B1*A1+B2*A7))/TEMP
VV = V*(1.+C4/TEMP)
RETURN
C IF TYPE=3 THE QUADRATIC IS ZEROED
30 UU = 0.D0
VV = 0.D0
RETURN
END
NEW 10
NEW 20
NEW 30
NEW 40
NEW 50
NEW 60
NEW 70
NEW 80
NEW 90
NEW 100
NEW 110
NEW 120
NEW 130
NEW 140
NEW 150
NEW 160
NEW 170
NEW 180
NEW 190
NEW 200
NEW 210
NEW 220
NEW 230
NEW 240
NEW 250
NEW 260
NEW 270
NEW 280
NEW 290
NEW 300
NEW 310
NEW 320
NEW 330
NEW 340
NEW 350
NEW 360
NEW 370
NEW 380
NEW 390
NEW 400

SUBROUTINE QUADSD(NN, U, V, P, Q, A, B)
C DIVIDES P BY THE QUADRATIC 1,U,V PLACING THE
C QUOTIENT IN Q AND THE REMAINDER IN A,B
DOUBLE PRECISION P(NN), Q(NN), U, V, A, B, C
INTEGER I
B = P(1)
Q(1) = B
A = P(2) - U*B
Q(2) = A
DO 10 I=3,NN
C = P(I) - U*A - V*B
QUA 10
QUA 20
QUA 30
QUA 40
QUA 50
QUA 60
QUA 70
QUA 80
QUA 90
QUA 100
QUA 110

```

```

      Q(I) = C
      B = A
      A = C
10  CONTINUE
      RETURN
      END
      SUBROUTINE QUAD(A, B1, C, SR, SI, LR, LI)
C  CALCULATE THE ZEROS OF THE QUADRATIC  $Az^2+B_1z+C$ .
C  THE QUADRATIC FORMULA, MODIFIED TO AVOID
C  OVERFLOW, IS USED TO FIND THE LARGER ZERO IF THE
C  ZEROS ARE REAL AND BOTH ZEROS ARE COMPLEX.
C  THE SMALLER REAL ZERO IS FOUND DIRECTLY FROM THE
C  PRODUCT OF THE ZEROS C/A.
      DOUBLE PRECISION A, B1, C, SR, SI, LR, LI, B,
      * D, E, DABS, DSQRT
      IF (A.NE.0.D0) GO TO 20
      SR = 0.D0
      IF (B1.NE.0.D0) SR = -C/B1
      LR = 0.D0
10  SI = 0.D0
      LI = 0.D0
      RETURN
20  IF (C.NE.0.D0) GO TO 30
      SR = 0.D0
      LR = -B1/A
      GO TO 10
C  COMPUTE DISCRIMINANT AVOIDING OVERFLOW
30  B = B1/2.D0
      IF (DABS(B).LT.DABS(C)) GO TO 40
      E = 1.D0 - (A/B)*(C/B)
      D = DSQRT(DABS(E))*DABS(B)
      GO TO 50
40  E = A
      IF (C.LT.0.D0) E = -A
      E = B*(B/DABS(C)) - E
      D = DSQRT(DABS(E))*DSQRT(DABS(C))
50  IF (E.LT.0.D0) GO TO 60
C  REAL ZEROS
      IF (B.GE.0.D0) D = -D
      LR = (-B+D)/A
      SR = 0.D0
      IF (LR.NE.0.D0) SR = (C/LR)/A
      GO TO 10
C  COMPLEX CONJUGATE ZEROS
60  SR = -B/A
      LR = SR
      SI = DABS(D/A)
      LI = -SI
      RETURN
      END

```

```

QUA 120
QUA 130
QUA 140
QUA 150
QUA 160
QUA 170
QUA 10
QUA 20
QUA 30
QUA 40
QUA 50
QUA 60
QUA 70
QUA 80
QUA 90
QUA 100
QUA 110
QUA 120
QUA 130
QUA 140
QUA 150
QUA 160
QUA 170
QUA 180
QUA 190
QUA 200
QUA 210
QUA 220
QUA 230
QUA 240
QUA 250
QUA 260
QUA 270
QUA 280
QUA 290
QUA 300
QUA 310
QUA 320
QUA 330
QUA 340
QUA 350
QUA 360
QUA 370
QUA 380
QUA 390
QUA 400
QUA 410
QUA 420
QUA 430
QUA 440

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ALGORITHM 494

PDEONE, Solutions of Systems of Partial Differential Equations [D3]

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Key Words and Phrases: partial differential equations, method of lines, ordinary differential equations

CR Categories: 5.17

Language: Fortran

DESCRIPTION

The description of this algorithm, test results, and references are contained in the authors' paper, "Software for Nonlinear Partial Differential Equations," *ACM Trans. Math. Software* 1, 3 (Sept. 1975), 232-260.

ALGORITHM

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SUBROUTINE PDEONE(T, U, UDOT, NPDE, NPTS)           PDE 10
  DIMENSION U(NPDE,NPTS), UDOT(NPDE,NPTS)         PDE 20
C PDEONE IS AN INTERFACE SUBROUTINE WHICH USES CENTERED DIF- PDE 30
C FERENCE APPROXIMATIONS TO CONVERT ONE-DIMENSIONAL SYSTEMS PDE 40
C OF PARTIAL DIFFERENTIAL EQUATIONS INTO A SYSTEM OF ORDINARY PDE 50
C DIFFERENTIAL EQUATIONS, UDOT = F(T,X,U). THIS ROUTINE IS PDE 60
C INTENDED TO BE USED WITH A ROBUST ODE INTEGRATOR. PDE 70
C INPUT.. PDE 80
C NPDE = NUMBER OF PARTIAL DIFFERENTIAL EQUATIONS. PDE 90
C NPTS = NUMBER OF SPATIAL GRID POINTS. PDE 100
C T = CURRENT VALUE OF TIME. PDE 110
C U = AN NPDE BY NPTS ARRAY CONTAINING THE COMPUTED PDE 120
C SOLUTION AT TIME T. PDE 130
C OUTPUT.. PDE 140
C UDOT = AN NPDE BY NPTS ARRAY CONTAINING THE RIGHT HAND PDE 150
C SIDE OF THE RESULTING SYSTEM OF ODE*S, F(T,X,U), PDE 160
C OBTAINED BY DISCRETIZING THE GIVEN PDE*S. PDE 170
C THE USER MUST INSERT A DIMENSION STATEMENT OF THE FOLLOW- PDE 180
C ING FORM.. PDE 190
C DIMENSION DVAL(**,**,2), UX(**), UAVG(**), ALPHA(**), PDE 200
C * BETA(**), GAMMA(**) PDE 210
C THE SYMBOLS ** DENOTE THE ACTUAL NUMERICAL VALUE OF NPDE PDE 220
C FOR THE PROBLEM BEING SOLVED. PDE 230
C COMMON BLOCK MESH CONTAINS THE USER SPECIFIED SPATIAL PDE 240
C GRID POINTS. PDE 250
C COMMON BLOCK COORD CONTAINS 0,1, OR 2 DEPENDING ON WHETHER PDE 260
C THE PROBLEM IS IN CARTESIAN, CYLINDRICAL, OR SPHERICAL PDE 270

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C COORDINATES, RESPECTIVELY. PDE 280
COMMON /MESH/ X(1) PDE 290
COMMON /COORD/ ICORD PDE 300
ICORD1 = ICORD + 1 PDE 310
C UPDATE THE LEFT BOUNDARY VALUES PDE 320
CALL BNDRY(T, X(1), U, ALPHA, BETA, GAMMA, NPDE) PDE 330
ITEST = 0 PDE 340
DXI = 1./((X(2)-X(1))) PDE 350
DO 10 K=1, NPDE PDE 360
IF (BETA(K).NE.0.0) GO TO 10 PDE 370
U(K,1) = GAMMA(K)/ALPHA(K) PDE 380
ITEST = ITEST + 1 PDE 390
10 CONTINUE PDE 400
IF (ITEST.EQ.NPDE) GO TO 50 PDE 410
IF (ITEST.EQ.0) GO TO 20 PDE 420
CALL BNDRY(T, X(1), U, ALPHA, BETA, GAMMA, NPDE) PDE 430
C EVALUATE D COEFFICIENTS AT THE LEFT BOUNDARY PDE 440
20 CALL D(T, X(1), U, DVAL, NPDE) PDE 450
C FORM APPROXIMATIONS TO DU/DX AT THE LEFT BOUNDARY PDE 460
DO 40 K=1, NPDE PDE 470
IF (BETA(K).NE.0.0) GO TO 30 PDE 480
UX(K) = DXI*(U(K,2)-U(K,1)) PDE 490
GO TO 40 PDE 500
30 UX(K) = (GAMMA(K)-ALPHA(K)*U(K,1))/BETA(K) PDE 510
40 CONTINUE PDE 520
C EVALUATE U-AVERAGE IN THE FIRST INTERVAL PDE 530
50 DO 60 K=1, NPDE PDE 540
UAVG(K) = .5*(U(K,2)+U(K,1)) PDE 550
60 CONTINUE PDE 560
C EVALUATE D COEFFICIENTS IN THE FIRST INTERVAL PDE 570
XAVGR = .5*(X(2)+X(1)) PDE 580
CALL D(T, XAVGR, UAVG, DVAL(1,1,2), NPDE) PDE 590
DXIL = 1. PDE 600
DXIR = DXI PDE 610
IF (ICORD.EQ.0) GO TO 70 PDE 620
DXIL = X(1)**ICORD PDE 630
DXIR = DXIR*XAVGR**ICORD PDE 640
C EVALUATE DUXX AT THE LEFT BOUNDARY PDE 650
70 IF (ITEST.EQ.NPDE) GO TO 100 PDE 660
DXIC = FLOAT(ICORD1)/(XAVGR**ICORD1-X(1)**ICORD1) PDE 670
DO 90 L=1, NPDE PDE 680
DO 80 K=1, NPDE PDE 690
DVAL(K,L,1) = DXIC*(DVAL(K,L,2)*(U(L,2)-U(L,1))* PDE 700
* DXIR-DVAL(K,L,1)*UX(L)*DXIL) PDE 710
80 CONTINUE PDE 720
90 CONTINUE PDE 730
C EVALUATE RIGHTHAND SIDE OF PDE* S AT THE LEFT BOUNDARY PDE 740
CALL F(T, X(1), U, UX, DVAL, UDOT, NPDE) PDE 750
C SET UDOT = 0 FOR KNOWN LEFT BOUNDARY VALUES PDE 760
100 DO 110 K=1, NPDE PDE 770
IF (BETA(K).EQ.0.0) UDOT(K,1) = 0.0 PDE 780
110 CONTINUE PDE 790
C UPDATE THE RIGHT BOUNDARY VALUES PDE 800
CALL BNDRY(T, X(NPTS), U(1,NPTS), ALPHA, BETA, GAMMA, NPDE) PDE 810
ITEST = 0 PDE 820
DO 120 K=1, NPDE PDE 830
IF (BETA(K).NE.0.0) GO TO 120 PDE 840
U(K,NPTS) = GAMMA(K)/ALPHA(K) PDE 850
ITEST = ITEST + 1 PDE 860
120 CONTINUE PDE 870
IBCK = 1 PDE 880
IFWD = 2 PDE 890
ILIM = NPTS - 1 PDE 900
C MAIN LOOP TO FORM ORDINARY DIFFERENTIAL EQUATIONS AT THE PDE 910
C GRID POINTS PDE 920
DO 160 I=2, ILIM PDE 930
K = IBCK PDE 940
IBCK = IFWD PDE 950
IFWD = K PDE 960
XAVGL = XAVGR PDE 970
XAVGR = .5*(X(I+1)+X(I)) PDE 980
DXI = 1./(X(I+1)-X(I-1)) PDE 990
DXIL = DXIR PDE 1000
DXIR = 1./(X(I+1)-X(I)) PDE 1010
IF (ICORD.NE.0) DXIR = DXIR*XAVGR**ICORD PDE 1020
DXIC = FLOAT(ICORD1)/(XAVGR**ICORD1-XAVGL**ICORD1) PDE 1030
C EVALUATE DU/DX AND U-AVERAGE AT THE I-TH GRID POINT AND PDE 1040
C INTERVAL RESPECTIVELY PDE 1050
DO 130 K=1, NPDE PDE 1060
UX(K) = DXI*(U(K,I+1)-U(K,I-1)) PDE 1070
UAVG(K) = .5*(U(K,I+1)+U(K,I)) PDE 1080
130 CONTINUE PDE 1090
C EVALUATE D COEFFICIENTS IN THE I-TH INTERVAL PDE 1100
CALL D(T, XAVGR, UAVG, DVAL(1,1,IFWD), NPDE) PDE 1110
C EVALUATE DUXX AT THE I-TH GRID POINT PDE 1120
DO 150 L=1, NPDE PDE 1130
DO 140 K=1, NPDE PDE 1140
DVAL(K,L,IBCK) = (DVAL(K,L,IFWD)*(U(L,I+1)-U(L,I))* PDE 1150
* DXIR-DVAL(K,L,IBCK)*(U(L,I)-U(L,I-1))*DXIL)*DXIC PDE 1160
140 CONTINUE PDE 1170
150 CONTINUE PDE 1180

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ALGORITHM 495

Solution of an Overdetermined System of Linear Equations in the Chebyshev Norm [F4]

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Key Words and Phrases: Chebyshev solution, overdetermined system, linear programming, simplex method

CR Categories: 3.15

Language: Fortran

DESCRIPTION

The algorithm calculates a Chebyshev (or l_∞) solution to an $m \times n$ overdetermined system of linear equations $Ax = b$, i.e. given equations

$$\sum_{j=1}^n a_{i,j} x_j = b_i, \quad i = 1, 2, \dots, m, \quad m \geq n,$$

the subroutine determines a vector $x^* = [x_1^*, x_2^*, \dots, x_n^*]^T$ which minimizes the maximum absolute value of the residuals

$$e(x) = \|b - Ax\|_\infty = \max_{1 \leq i \leq m} \left| b_i - \sum_{j=1}^n a_{i,j} x_j \right|. \quad (1)$$

The algorithm can be used to solve the linear Chebyshev data fitting problem. Suppose that data consisting of m points (z_i, y_i) are to be approximated by a linear approximating function $\alpha_1 \phi_1(z) + \alpha_2 \phi_2(z) + \dots + \alpha_n \phi_n(z)$ in the Chebyshev norm. This is equivalent to finding a Chebyshev solution to the system of linear equations

$$\sum_{j=1}^n \phi_j(z_i) \alpha_j = y_i, \quad i = 1, 2, \dots, m.$$

If the y_i values contain only small inherent errors, a Chebyshev approximation may be preferred to a least-squares or least-first-power approximation.

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The algorithm is a modification of the simplex method of linear programming applied to the *dual* formulation of the Chebyshev problem. It does not require that the $m \times n$ matrix of coefficients $A = \{a_{i,j}\}$ satisfy the Haar condition, nor does it require that it be of full rank. Computational experience with this and other algorithms indicates that it is efficient in both time and storage requirements. An approximate solution is determined during the first $k + 1$ iterations, where k denotes the rank of the matrix A , and in the subsequent iterations this solution is refined until either (i) a Chebyshev solution x^* is obtained which minimizes $e(x)$, or (ii) an approximate solution \hat{x} is found for which

$$(e(\hat{x}) - e(x^*)) / e(x^*) < RELERR \quad (2)$$

where *RELERR* is a real variable specified by the user (see below).

The parameters M and N represent the number of equations and number of unknowns, respectively. *MDIM* and *NDIM* should be set to at least $M + 1$ and $N + 3$, respectively. The simplex iterations are carried out in the two-dimensional real array A of size $(NDIM, MDIM)$. Initially, the *transpose* of the matrix A must be stored in the first N rows and M columns of A ; this arrangement allows the code to correspond closely to the formal description of the algorithm [1]. The right-hand-side vector b must be stored in the array B . These initial coefficients assigned to A and B are subsequently destroyed. *TOL* is a real variable which should be set to a small positive value; e.g. $TOL = 10^{-d+1}$ where d represents the number of decimal digits of accuracy available. (On an IBM 370 we usually set *TOL* equal to 10^{-6} when using single precision arithmetic and equal to 10^{-15} when using double precision arithmetic.) Essentially, *the subroutine cannot distinguish between zero and any quantity whose magnitude does not exceed TOL*. In particular, it will not pivot on any number whose magnitude does not exceed *TOL*. (Consequently, the subroutine may terminate prematurely if the system of equations is poorly scaled.) The parameter *RELERR* should be set to 0.0 if a Chebyshev solution x^* is required. If *RELERR* is set to a positive value, the subroutine determines an approximate solution \hat{x} satisfying inequality (2); the usual effect of this option (say, with *RELERR* set to 0.1) is a saving in the number of simplex iterations required.

On exit, if *RELERR* was set to 0.0, the array X contains a Chebyshev solution x^* . However, if the initial value of *RELERR* was positive, then X contains an approximate solution \hat{x} . The subroutine assigns a new value to *RELERR*, smaller than its initial value, and \hat{x} satisfies (2) for this new value of *RELERR*. The rank of the $m \times n$ matrix A is determined during the first k iterations and assigned to the integer *RANK*. (The rank estimate *RANK* is dependent upon *TOL*; decreasing the value of *TOL* may increase the value of *RANK*.) The residuals

$$b_i - \sum_{j=1}^n a_{i,j} x_j, \quad i = 1, 2, \dots, m \quad (3)$$

are assigned in order to the array B . *RESMAX* supplies the absolute value of the largest residual; this is the minimum value of (1) if *RELERR* was set to 0.0. If *RELERR* was set to a positive value, then *RESMAX* gives the value of (1) corresponding to the approximate solution \hat{x} . (The subroutine, which overwrites the original data A and b , supplies the residuals automatically. However, these automatic residuals may contain few significant figures, especially when *RESMAX* is within one or two orders of magnitude of *TOL*. Indeed, if $RESMAX \leq TOL$ the automatic residuals may all be set to zero by the subroutine. *It is therefore often advisable to obtain the residuals directly from (3) as an additional calculation.*) *ITER* records the number of iterations performed by the simplex method. Finally, *OCODE* is an exit code with the value 1 if a solution is calculated successfully, and 2 if the calculations are terminated prematurely. This latter condition occurs only when a pivot is encountered whose magnitude does not exceed *TOL*, and in this event all output information pertains to the last completed simplex iteration. Since a

Chebyshev solution is not necessarily unique, the subroutine attempts to determine if other optimal solutions exist. An exit code of 1 indicates that the solution is unique, while an exit code of 0 indicates that the solution is almost certainly not unique. (This uncertainty can only be resolved by a close examination of the final simplex tableau contained in A , and we do not consider such an examination to be warranted in practice). A Chebyshev solution may be nonunique simply because the matrix of coefficients A is not of full rank.

Example. Approximate $y(z) = e^z$ by $\alpha_1 + \alpha_2 z + \alpha_3 z^2 + \alpha_4 z^3$ in the Chebyshev norm on the 51 points defined by $z = 0(0.02)1$. Equivalently, calculate the Chebyshev solution to the system of equations $\sum_{j=1}^4 z_i^{j-1} x_j = e^{z_i}$, $i = 1, 2, \dots, 51$, where $z_i = 0.02(i - 1)$. Setting TOL equal to 10^{-6} and setting $RELERR$ to 0.0, the subroutine gives the following solution on an IBM 370/145 using single precision arithmetic (approximately 7 decimal digits): $x_1^* = 0.999456$, $x_2^* = 1.016599$, $x_3^* = 0.421690$, $x_4^* = 0.279989$. The number of iterations required is 8, and the value of $RESMAX$ is 0.000544.

REFERENCES

1. BARRODALE, I., AND PHILLIPS, C. An improved algorithm for discrete Chebyshev linear approximations. Proc. 4th Manitoba Conf. on Numerical Mathematics, U. of Manitoba, Winnipeg, Canada, 1974, pp. 177-190.

ALGORITHM

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SUBROUTINE CHEB(M, N, MDIM, NDIM, A, B, TOL, RELERR, X, RANK,      CHE 10
* RESMAX, ITER, OCODE)                                          CHE 20
C THIS SUBROUTINE USES A MODIFICATION OF THE SIMPLEX METHOD      CHE 30
C OF LINEAR PROGRAMMING TO CALCULATE A CHEBYSHEV SOLUTION TO   CHE 40
C AN OVER-DETERMINED SYSTEM OF LINEAR EQUATIONS.              CHE 50
C DESCRIPTION OF PARAMETERS.                                     CHE 60
C M      NUMBER OF EQUATIONS.                                    CHE 70
C N      NUMBER OF UNKNOWNNS (N MUST NOT EXCEED M).           CHE 80
C MDIM   THE NUMBER OF COLUMNS OF A, AT LEAST M+1.           CHE 90
C NDIM   THE NUMBER OF ROWS OF A, AT LEAST N+3.                CHE 100
C A      TWO DIMENSIONAL REAL ARRAY OF SIZE (NDIM,MDIM).      CHE 110
C        ON ENTRY, THE TRANSPOSE OF THE MATRIX OF              CHE 120
C        COEFFICIENTS OF THE OVER-DETERMINED SYSTEM MUST     CHE 130
C        BE STORED IN THE FIRST M COLUMNS AND N ROWS OF A.   CHE 140
C        THESE VALUES ARE DESTROYED BY THE SUBROUTINE.       CHE 150
C B      ONE DIMENSIONAL REAL ARRAY OF SIZE MDIM. ON ENTRY,   CHE 160
C        B MUST CONTAIN THE RIGHT-HAND SIDES OF THE           CHE 170
C        EQUATIONS IN ITS FIRST M LOCATIONS. ON EXIT, B        CHE 180
C        CONTAINS THE RESIDUALS FOR THE EQUATIONS IN ITS      CHE 190
C        FIRST M LOCATIONS (SEE DESCRIPTION).                  CHE 200
C TOL    A SMALL POSITIVE TOLERANCE. EMPIRICAL EVIDENCE      CHE 210
C        SUGGESTS TOL=10**(-D+1) WHERE D REPRESENTS THE      CHE 220
C        NUMBER OF DECIMAL DIGITS OF ACCURACY AVAILABLE       CHE 230
C        (SEE DESCRIPTION).                                    CHE 240
C RELERR A REAL VARIABLE WHICH ON ENTRY MUST HAVE THE VALUE   CHE 250
C        0.0 IF A CHEBYSHEV SOLUTION IS REQUIRED. IF RELERR    CHE 260
C        IS POSITIVE, THE SUBROUTINE CALCULATES AN            CHE 270
C        APPROXIMATE SOLUTION WITH RELERR AS AN UPPER BOUND   CHE 280
C        ON THE RELATIVE ERROR OF ITS LARGEST RESIDUAL (SEE   CHE 290
C        DESCRIPTION-INEQUALITY (2)). ON EXIT, THE VALUE OF    CHE 300
C        RELERR GIVES A SMALLER UPPER BOUND FOR THIS         CHE 310
C        RELATIVE ERROR.                                       CHE 320
C X      ONE DIMENSIONAL REAL ARRAY OF SIZE NDIM. ON EXIT,    CHE 330
C        THIS ARRAY CONTAINS A SOLUTION TO THE PROBLEM IN     CHE 340
C        ITS FIRST N LOCATIONS.                                 CHE 350
C RANK   AN INTEGER WHICH GIVES ON EXIT THE RANK OF THE       CHE 360
C        MATRIX OF COEFFICIENTS.                               CHE 370
C RESMAX THE LARGEST RESIDUAL IN MAGNITUDE.                    CHE 380
C ITER   THE NUMBER OF SIMPLEX ITERATIONS PERFORMED.         CHE 390
C OCODE  AN EXIT CODE WITH VALUES..                           CHE 400
C        0 - OPTIMAL SOLUTION WHICH IS PROBABLY                CHE 410
C          NON-UNIQUE (SEE DESCRIPTION).                       CHE 420
C        1 - UNIQUE OPTIMAL SOLUTION.                          CHE 430
C        2 - CALCULATIONS TERMINATED PREMATURELY              CHE 440
C          DUE TO ROUNDING ERRORS.                             CHE 450
C IF YOUR FORTRAN COMPILER PERMITS A SINGLE COLUMN OF A TWO   CHE 460
C DIMENSIONAL ARRAY TO BE PASSED TO A ONE DIMENSIONAL ARRAY CHE 470
C THROUGH A SUBROUTINE CALL, CONSIDERABLE SAVINGS IN         CHE 480
C EXECUTION TIME MAY BE ACHIEVED THROUGH THE USE OF THE      CHE 490
C FOLLOWING SUBROUTINE WHICH OPERATES ON COLUMN VECTORS.     CHE 500
C   SUBROUTINE COL (V1,V2,MLT,NOTROW,I1,NP2)                   CHE 510
C THIS SUBROUTINE SUBTRACTS FROM THE VECTOR V1 A MULTIPLE OF CHE 520
C THE VECTOR V2 STARTING AT THE I1*TH ELEMENT UP TO THE     CHE 530
C NP2*TH ELEMENT, EXCEPT FOR THE NOTROW*TH ELEMENT.        CHE 540
C   REAL V1(NP2),V2(NP2),MLT                                  CHE 550
C   DO 1 I=I1,NP2                                           CHE 560

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C      IF(I.EQ.NOTROW) GO TO 1          CHE 570
C      V1(I)=V1(I)-MLT*V2(I)           CHE 580
C 1    CONTINUE                        CHE 590
C      RETURN                          CHE 600
C      END                              CHE 610
C SEE COMMENTS FOLLOWING STATEMENT NUMBER 340 FOR          CHE 620
C INSTRUCTIONS ON THE IMPLEMENTATION OF THIS MODIFICATION. CHE 630
      DIMENSION A(NDIM,MDIM), B(MDIM), X(NDIM)             CHE 640
      INTEGER PROW, PCOL, RANK, RANKP1, OCODE             CHE 650
C BIG MUST BE SET EQUAL TO ANY VERY LARGE REAL CONSTANT. CHE 660
C ITS VALUE HERE IS APPROPRIATE FOR THE IBM/370/145.     CHE 670
      DATA BIG /1.E+75/                                  CHE 680
C INITIALIZATION.                                        CHE 690
      MP1 = M + 1                                         CHE 700
      NP1 = N + 1                                         CHE 710
      NP2 = N + 2                                         CHE 720
      NP3 = N + 3                                         CHE 730
      NP1MR = 1                                           CHE 740
      RANK = N                                             CHE 750
      RELTMP = RELERR                                     CHE 760
      RELERR = 0.                                         CHE 770
      DO 10 J=1,M                                         CHE 780
        A(NP1,J) = 1.                                     CHE 790
        A(NP2,J) = -B(J)                                 CHE 800
        A(NP3,J) = N + J                                 CHE 810
10    CONTINUE                                           CHE 820
      A(NP1,MP1) = 0.                                     CHE 830
      ITER = 0                                            CHE 840
      OCODE = 1                                          CHE 850
      DO 20 I=1,N                                         CHE 860
        X(I) = 0.                                        CHE 870
        A(I,MP1) = I                                    CHE 880
20    CONTINUE                                           CHE 890
C LEVEL 1.                                              CHE 900
      LEV = 1                                           CHE 910
      K = 0                                              CHE 920
30    K = K + 1                                          CHE 930
      KP1 = K + 1                                        CHE 940
      NP1MK = NP1 - K                                   CHE 950
      MODE = 0                                          CHE 960
      DO 40 J=K,M                                        CHE 970
        B(J) = 1.                                       CHE 980
40    CONTINUE                                           CHE 990
C DETERMINE THE VECTOR TO ENTER THE BASIS.             CHE 1000
50    D = -BIG                                          CHE 1010
      DO 60 J=K,M                                        CHE 1020
        IF (B(J).EQ.0.) GO TO 60                       CHE 1030
        DD = ABS(A(NP2,J))                              CHE 1040
        IF (DD.LE.D) GO TO 60                          CHE 1050
        PCOL = J                                        CHE 1060
        D = DD                                          CHE 1070
60    CONTINUE                                           CHE 1080
      IF (K.GT.1) GO TO 70                              CHE 1090
C TEST FOR ZERO RIGHT-HAND SIDE.                       CHE 1100
      IF (D.GT.TOL) GO TO 70                          CHE 1110
      RESMAX = 0.                                       CHE 1120
      MODE = 2                                          CHE 1130
      GO TO 380                                         CHE 1140
C DETERMINE THE VECTOR TO LEAVE THE BASIS.             CHE 1150
70    D = TOL                                          CHE 1160
      DO 80 I=1,NP1MK                                   CHE 1170
        DD = ABS(A(I,PCOL))                             CHE 1180
        IF (DD.LE.D) GO TO 80                          CHE 1190
        PROW = I                                       CHE 1200
        D = DD                                          CHE 1210
80    CONTINUE                                           CHE 1220
      IF (D.GT.TOL) GO TO 330                           CHE 1230
C CHECK FOR LINEAR DEPENDENCE IN LEVEL 1.              CHE 1240
      B(PCOL) = 0.                                     CHE 1250
      IF (MODE.EQ.1) GO TO 50                          CHE 1260
      DO 100 J=K,M                                      CHE 1270
        IF (B(J).EQ.0.) GO TO 100                     CHE 1280
        DO 90 I=1,NP1MK                                CHE 1290
          IF (ABS(A(I,J)).LE.TOL) GO TO 90             CHE 1300
          MODE = 1                                     CHE 1310
          GO TO 50                                     CHE 1320
90    CONTINUE                                           CHE 1330
100   CONTINUE                                           CHE 1340
      RANK = K - 1                                     CHE 1350
      NP1MR = NP1 - RANK                               CHE 1360
      OCODE = 0                                        CHE 1370
      GO TO 160                                         CHE 1380
110   IF (PCOL.EQ.K) GO TO 130                         CHE 1390
C INTERCHANGE COLUMNS IN LEVEL 1.                    CHE 1400
      DO 120 I=1,NP3                                    CHE 1410
        D = A(I,PCOL)                                  CHE 1420
        A(I,PCOL) = A(I,K)                            CHE 1430
        A(I,K) = D                                    CHE 1440
120   CONTINUE                                           CHE 1450
130   IF (PROW.EQ.NP1MK) GO TO 150                   CHE 1460
C INTERCHANGE ROWS IN LEVEL 1.                        CHE 1470
      DO 140 J=1,MP1                                    CHE 1480
        D = A(PROW,J)                                  CHE 1490

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        A(PROW,J) = A(NP1MK,J)
        A(NP1MK,J) = D
140 CONTINUE
150 IF (K.LT.N) GO TO 30
160 IF (RANK.EQ.M) GO TO 380
    RANKP1 = RANK + 1
C LEVEL 2.
    LEV = 2
C DETERMINE THE VECTOR TO ENTER THE BASIS
    D = TOL
    DO 170 J=RANKP1,M
        DD = ABS(A(NP2,J))
        IF (DD.LE.D) GO TO 170
        PCOL = J
        D = DD
170 CONTINUE
C COMPARE CHEBYSHEV ERROR WITH TOL.
    IF (D.GT.TOL) GO TO 180
    RESMAX = 0.
    MODE = 3
    GO TO 380
180 IF (A(NP2,PCOL).LT.-TOL) GO TO 200
    A(NP1,PCOL) = 2. - A(NP1,PCOL)
    DO 190 I=NP1MR,NP3
        IF (I.EQ.NP1) GO TO 190
        A(I,PCOL) = -A(I,PCOL)
190 CONTINUE
C ARRANGE FOR ALL ENTRIES IN PIVOT COLUMN
C (EXCEPT PIVOT) TO BE NEGATIVE.
200 DO 220 I=NP1MR,N
    IF (A(I,PCOL).LT.TOL) GO TO 220
    DO 210 J=1,M
        A(NP1,J) = A(NP1,J) + 2.*A(I,J)
        A(I,J) = -A(I,J)
210 CONTINUE
    A(I,MP1) = -A(I,MP1)
220 CONTINUE
    PROW = NP1
    GO TO 330
230 IF (RANKP1.EQ.M) GO TO 380
    IF (PCOL.EQ.M) GO TO 250
C INTERCHANGE COLUMNS IN LEVEL 2.
    DO 240 I=NP1MR,NP3
        D = A(I,PCOL)
        A(I,PCOL) = A(I,M)
        A(I,M) = D
240 CONTINUE
250 MM1 = M - 1
C LEVEL 3.
    LEV = 3
C DETERMINE THE VECTOR TO ENTER THE BASIS.
260 D = -TOL
    VAL = 2.*A(NP2,M)
    DO 280 J=RANKP1,MM1
        IF (A(NP2,J).GE.D) GO TO 270
        PCOL = J
        D = A(NP2,J)
        MODE = 0
        GO TO 280
270 DD = VAL - A(NP2,J)
    IF (DD.GE.D) GO TO 280
    MODE = 1
    PCOL = J
    D = DD
280 CONTINUE
    IF (D.GE.-TOL) GO TO 380
    DD = -D/A(NP2,M)
    IF (DD.GE.RELTMP) GO TO 290
    RELERR = DD
    MODE = 4
    GO TO 380
290 IF (MODE.EQ.0) GO TO 310
    DO 300 I=NP1MR,NP1
        A(I,PCOL) = 2.*A(I,M) - A(I,PCOL)
300 CONTINUE
    A(NP2,PCOL) = D
    A(NP3,PCOL) = -A(NP3,PCOL)
C DETERMINE THE VECTOR TO LEAVE THE BASIS.
310 D = BIG
    DO 320 I=NP1MR,NP1
        IF (A(I,PCOL).LE.TOL) GO TO 320
        DD = A(I,M)/A(I,PCOL)
        IF (DD.GE.D) GO TO 320
        PROW = I
        D = DD
320 CONTINUE
    IF (D.LT.BIG) GO TO 330
    OCODE = 2
    GO TO 380
C PIVOT ON A(PROW,PCOL).
330 PIVOT = A(PROW,PCOL)
    DO 340 J=1,M
        A(PROW,J) = A(PROW,J)/PIVOT

```

CHE 1500
CHE 1510
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CHE 2370
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CHE 2390
CHE 2400
CHE 2410
CHE 2420

```

340 CONTINUE
C IF PERMITTED, USE SUBROUTINE COL IN THE DESCRIPTION
C SECTION AND REPLACE THE FOLLOWING EIGHT STATEMENTS DOWN TO
C AND INCLUDING STATEMENT NUMBER 360 BY..
C DO 360 J=1,M
C IF (J.EQ.PCOL) GO TO 360
C CALL COL (A(1,J),A(1,PCOL),A(1,PCOL),PROW,PROW,NP1MR,NP2)
C 360 CONTINUE
DO 360 J=1,M
IF (J.EQ.PCOL) GO TO 360
D = A(1,PCOL)
DO 350 I=NP1MR,NP2
IF (I.EQ.PROW) GO TO 350
A(I,J) = A(I,J) - D*A(I,PCOL)
350 CONTINUE
360 CONTINUE
TPIVOT = -PIVOT
DO 370 I=NP1MR,NP2
A(I,PCOL) = A(I,PCOL)/TPIVOT
370 CONTINUE
A(1,PCOL) = 1./PIVOT
D = A(1,MP1)
A(1,MP1) = A(NP3,PCOL)
A(NP3,PCOL) = D
ITER = ITER + 1
GO TO (110, 230, 260), LEV
C PREPARE OUTPUT.
380 DO 390 J=1,M
B(J) = 0.
390 CONTINUE
IF (MODE.EQ.2) GO TO 450
DO 400 J=1,RANK
K = A(NP3,J)
X(K) = A(NP2,J)
400 CONTINUE
IF (MODE.EQ.3 .OR. RANK.EQ.M) GO TO 450
DO 410 I=NP1MR,NP1
K = ABS(A(I,MP1)) - FLOAT(N)
B(K) = A(NP2,M)*SIGN(1.,A(I,MP1))
410 CONTINUE
IF (RANKP1.EQ.M) GO TO 430
DO 420 J=RANKP1,MM1
K = ABS(A(NP3,J)) - FLOAT(N)
B(K) = (A(NP2,M)-A(NP2,J))*SIGN(1.,A(NP3,J))
420 CONTINUE
C TEST FOR NON-UNIQUE SOLUTION.
430 DO 440 I=NP1MR,NP1
IF (ABS(A(I,M)).GT.TOL) GO TO 440
OCCODE = 0
GO TO 450
440 CONTINUE
450 IF (MODE.NE.2 .AND. MODE.NE.3) RESMAX = A(NP2,M)
IF (RANK.EQ.M) RESMAX = 0.
IF (MODE.EQ.4) RESMAX = RESMAX - D
RETURN
END
CHE 2430
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CHE 2880
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ALGORITHM 496

The LZ Algorithm To Solve the Generalized Eigenvalue Problem for Complex Matrices [F2]

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Key Words and Phrases: eigenvalues, generalized eigenvalue problem

CR Categories: 5.14

Language: Fortran

DESCRIPTION

The program given here is an implementation of the LZ algorithm [1] for finding \mathbf{x} and λ such that $A\mathbf{x} = \lambda B\mathbf{x}$, where A and B are $n \times n$ matrices. The LZ algorithm does not require matrix inversion and may be used when either or both matrices are singular. It is a generalization of Rutishauser's LR method [3] for the standard eigenvalue problem $A\mathbf{x} = \lambda\mathbf{x}$ and closely resembles the QZ algorithm given by Moler and Stewart [2] for the generalized problem given above. Since the LZ algorithm uses elementary transformations and the QZ algorithm uses orthogonal transformations, the LZ algorithm is probably more efficient when either A or B is complex.

The LZ algorithm is based on three observations:

(1) If L and M are nonsingular matrices, the eigenvalue problems $LAM\mathbf{y} = \lambda LBM\mathbf{y}$ and $A\mathbf{x} = \lambda B\mathbf{x}$ have the same eigenvalues and their eigenvectors are related by $\mathbf{x} = M\mathbf{y}$.

(2) If A is a triangular matrix with diagonal elements α_i , and B is a triangular matrix with diagonal elements β_i , then for each i , $i = 1, 2, \dots, n$, α_i/β_i is an eigenvalue of the generalized problem if $\beta_i \neq 0$.

(3) There exist matrices L and M such that LAM and LBM are upper triangular and L and M are the products of lower triangular and permutation matrices.

The aim of the LZ algorithm is to find the matrices L and M which reduce A and B simultaneously to triangular form. It accomplishes this aim in two phases.

(1) The first phase simultaneously reduces A to upper Hessenberg form (i.e. $a_{i,j} = 0$ for $i > j$) and B to upper triangular form by (a) reducing B to triangular form as in Gaussian elimination and applying the transformations to A , and (b) reducing A to upper Hessenberg form while preserving the triangularity of B by using row transformations to zero the elements below the subdiagonal of A and column transformations to zero the subdiagonal elements of B introduced by the row transformations.

(2) The second phase iteratively reduces A to upper triangular form while preserving the triangularity of B where each iteration is given by (a) finding a shift λ_k , (b) finding a stabilized transformation \bar{L}_k such that $\bar{L}_k(A_k - \lambda_k B_k)$ is

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upper triangular, (c) finding a stabilized transformation \bar{M}_k such that $\bar{L}_k B_k \bar{M}_k$ is upper triangular, (d) setting $A_{k+1} = \bar{L}_k A_k \bar{M}_k$ and $B_{k+1} = \bar{L}_k B_k \bar{M}_k$.

If \bar{L}_k and \bar{M}_k are constructed appropriately, A_{k+1} will be upper Hessenberg and hopefully some of its subdiagonal elements will be small enough to be considered zero.

The shift λ_k is used to hasten the convergence of the algorithm. In practice it is set to the solution of the lowest 2×2 subproblem on the diagonal of $A_k - \lambda B_k$ which has not been triangularized.

Each matrix \bar{L}_k is a product of matrices $L_{k,n-1} L_{k,n-2} \dots L_{k,2} L_{k,1}$ and each \bar{M}_k is a product of matrices $M_{k,1} M_{k,2} \dots M_{k,n-1}$, where each $L_{k,i}$ and each $M_{k,i}$ is a stabilized elementary transformation (see [4]). The matrix $L_{k,1}$ is designed to annihilate the (2, 1) element of $A_k - \lambda_k B_k$. For $i \geq 1$, $M_{k,i}$ zeros the i th subdiagonal element of B introduced by $L_{k,i}$, and for $i > 1$, $L_{k,i}$ zeros the $(i + 1, i - 1)$ -th element of A introduced by $M_{k,i-1}$.

The LZ algorithm determines the eigenvectors of the problem by computing the eigenvectors of the triangularized A and B and multiplying them by the product of all the column transformations.

The Fortran program given here is designed to find the eigensystem for two complex matrices A and B and consists of two subroutines which must be called separately. A user should first issue a call to LZHES to reduce A to upper Hessenberg form and B to triangular form and then issue a call to LZHES to determine the eigensystem of the reduced matrices. If A and B have already been reduced to Hessenberg and triangular form, then calling LZHES is unnecessary. In this situation, if eigenvectors are requested, the matrix X should be set to the identity matrix. The subroutine LZIT does not return the eigenvalues of the problem but the diagonal elements of the triangularized matrices. The i th eigenvalue may be found by dividing EIGA(i) by EIGB(i). The user is warned that EIGB(i) may be zero. The eigenvectors produced by LZIT are normalized so that the modulus of their largest component is 1.

Entrance to LZHES is attained by the statement

```
CALL LZHES(N,A,NA,B,NB,X,NX,WANTX)
```

where the input parameters are

N	order of the A and B matrices,
A	$n \times n$ complex matrix,
NA	row dimension of A,
B	$n \times n$ complex matrix,
NB	row dimension of B,
NX	row dimension of X (see output parameters)
WANTX	logical variable which should be set to .TRUE. if eigenvectors are wanted and otherwise set to .FALSE.,

and the output parameters are

A	complex upper Hessenberg matrix; the original A matrix is destroyed,
B	complex upper triangular matrix; the original B matrix is destroyed,
X	$n \times n$ complex array containing the transformations needed to compute the eigenvectors of the problem if WANTX was .TRUE. .

Entrance to LZIT is attained by the statement

```
CALL LZIT(N,A,NA,B,NB,X,NX,WANTX,ITER,EIGA,EIGB)
```

where the input parameters are

N	order of the A and B matrices,
A	$n \times n$ complex upper Hessenberg matrix,
NA	row dimension of the A matrix,
B	$n \times n$ complex upper triangular matrix,
NB	row dimension of the B matrix,
X	$n \times n$ complex array containing the transformations to determine the

eigenvectors of the original problem if WANTX has been set to .TRUE. ,
 NX row dimension of the X matrix,
 WANTX logical variable which should be set to .TRUE. if eigenvectors are wanted and otherwise set to .FALSE. ,

and the output parameters are

X $n \times n$ complex array whose i th column contains the i th eigenvectors if WANTX was set to .TRUE. ,
 ITER integer array of length n whose i th entry contains the number of iterations needed to find the i th eigenvalue; for any i , if $iter(i) = -1$, then after 30 iterations there was not a sufficient decrease in the last subdiagonal element of A to continue iterating so that not all the eigenvalues have been isolated,
 EIGA complex array of length n containing the diagonal of the triangularized A,
 EIGB complex array of length n containing the diagonal of B.

Since $|x|$, where x is a complex scalar, must be computed at least $8n$ times per iteration, the execution time of the method is dependent on the algorithm used to compute $|x|$. Replacing the true modulus by an approximation has little effect on the numerical properties but may result in a significantly more efficient code. Several versions of the LZ algorithm with different methods for computing $|x|$ were run on the IBM 360 model 67 with the Fortran H compiler $opt=2$. The results were interesting:

(1) The procedure CABS supplied by the manufacturer to compute $|x|$ was by far the slowest.

THE MATRIX A:

-238+	-344I	86+	178I	164+	240I	-166+	-308I	56+	158I
76+	152I	-96+	-128I	40+	-32I	60+	184I	-60+	-136I
118+	284I	55+	-182I	-13+	460I	34+	-192I	-176+	-214I
-314+	-160I	132+	78I	114+	296I	-90+	-164I	-424+	-374I
-54+	-24I	-205+	-400I	109+	148I	158+	312I	-38+	-96I

THE MATRIX B:

388+	94I	-386+	-122I	-250+	-14I	556+	130I	-396+	-62I
-304+	-76I	384+	64I	-160+	16I	-240+	-92I	240+	68I
-658+	-136I	-73+	100I	-109+	-250I	-118+	100I	406+	96I
-640+	-10I	204+	-42I	-692+	-90I	288+	66I	-172+	154I
-162+	-72I	631+	158I	131+	52I	-758+	-184I	278+	76I

TRUE EIGENVALUE		COMPUTED EIGENVALUE		RELATIVE ERROR		RELATIVE RESIDUAL		NO. OF ITERATIONS	
1	7.6470588235294E-01+	9.4117647058823E-01 I	7.6470588235296E-01+	9.4117647058821E-01 I	2.7009971587330E-14	4.3138589276375E-15	0		
2	-1.0000000000000E+00+	-1.3333333333333E+00 I	-1.0000000000000E+00+	-1.3333333333333E+00 I	6.0291550413457E-15	6.7886927362343E-15	1		
3	-3.5294117647059E-01+	-4.1176470588235E-01 I	-3.5294117647059E-01+	-4.1176470588233E-01 I	3.9441522601135E-14	2.9742630452889E-15	0		
4	-3.5294117647059E-01+	-4.1176470588235E-01 I	-3.5294117647058E-01+	-4.1176470588236E-01 I	1.1809767484717E-14	6.0057593010225E-15	3		
5	-3.5294117647059E-01+	4.1176470588235E-01 I	-3.5294117647058E-01+	4.1176470588232E-01 I	5.6542864065059E-14	8.3080870692588E-15	1		

Fig. 1

(2) Substituting the procedure RABS given by

```

REAL FUNCTION RABS(X)
COMPLEX X, XX
REAL T(2)
EQUIVALENCE (XX, T(1))
XX = X
RABS = ABS(T(1)) + ABS(T(2))
RETURN
END

```

for CABS decreased the time by about 20 percent when no eigenvectors were computed. When using RABS, changing the precision of the program requires very little effort.

(3) Computing $|x|$ as $\text{ABS}(\text{REAL}(X)) + \text{ABS}(\text{AIMAG}(X))$ as in the program given here decreased execution times from version 2 by roughly 10 percent when no eigenvectors were computed.

Although theoretically numerical growth can be quite large, our experiments have shown that the algorithm behaves like most other methods which use stabilized elementary transformations: numerical growth rarely occurs and errors rarely accumulate. The eigenvalues are usually determined to the accuracy justified by the condition of the problem.

The example shown in Figure 1 was generated in integer arithmetic by multiplying two diagonal matrices by nonsingular transformations. It was run on the CDC Run compiler. The relative residual is the quantity,

$$\|\beta_i A x_i - \alpha_i B x_i\|_\infty / (\|\beta_i\| \|A\|_\infty + \|\alpha_i\| \|B\|_\infty)$$

where α_i and β_i are the i th diagonal elements of the triangularized A and B matrices and x_i is the i th eigenvector.

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ALGORITHM

```

SUBROUTINE LZHS(N, A, NA, B, NB, X, NX, WANTX)           LZH 10
C THIS SUBROUTINE REDUCES THE COMPLEX MATRIX A TO UPPER LZH 20
C HESSENBERG FORM AND REDUCES THE COMPLEX MATRIX B TO LZH 30
C TRIANGULAR FORM                                     LZH 40
C INPUT PARAMETERS..                                  LZH 50
C N THE ORDER OF THE A AND B MATRICES                LZH 60
C A A COMPLEX MATRIX                                  LZH 70
C NA THE ROW DIMENSION OF THE A MATRIX               LZH 80
C B A COMPLEX MATRIX                                  LZH 90
C NB THE ROW DIMENSION OF THE B MATRIX               LZH 100
C NX THE ROW DIMENSION OF THE X MATRIX              LZH 110
C WANTX A LOGICAL VARIABLE WHICH IS SET TO .TRUE. IF LZH 120
C THE EIGENVECTORS ARE WANTED. OTHERWISE IT SHOULD LZH 130
C BE SET TO .FALSE.                                  LZH 140
C OUTPUT PARAMETERS..                                  LZH 150
C A ON OUTPUT A IS AN UPPER HESSENBERG MATRIX, THE LZH 160
C ORIGINAL MATRIX HAS BEEN DESTROYED                 LZH 170
C B AN UPPER TRIANGULAR MATRIX, THE ORIGINAL MATRIX LZH 180
C HAS BEEN DESTROYED                                 LZH 190
C X CONTAINS THE TRANSFORMATIONS NEEDED TO COMPUTE LZH 200
C THE EIGENVECTORS OF THE ORIGINAL SYSTEM           LZH 210
C COMPLEX Y, W, Z, A(NA,N), B(NB,N), X(NX,N)       LZH 220
C REAL C, D                                           LZH 230
C LOGICAL WANTX                                       LZH 240
C NM1 = N - 1                                         LZH 250
C REDUCE B TO TRIANGULAR FORM USING ELEMENTARY      LZH 260
C TRANSFORMATIONS                                     LZH 270
C DO 80 I=1,NM1                                       LZH 280
C D = 0.00                                            LZH 290
C IPI = I + 1                                         LZH 300

```

```

DO 10 K=IPL,N
  Y = B(K,I)
  C = ABS(REAL(Y)) + ABS(AIMAG(Y))
  IF (C.LE.D) GO TO 10
  D = C
  II = K
10  CONTINUE
  IF (D.EQ.0.0) GO TO 80
  Y = B(I,I)
  IF (D.LE.ABS(REAL(Y))+ABS(AIMAG(Y))) GO TO 40
C MUST INTERCHANGE
  DO 20 J=1,N
    Y = A(I,J)
    A(I,J) = A(II,J)
    A(II,J) = Y
  20  CONTINUE
  DO 30 J=I,N
    Y = B(I,J)
    B(I,J) = B(II,J)
    B(II,J) = Y
  30  CONTINUE
  DO 40 J=IPL,N
    Y = B(J,I)/B(I,I)
    IF (REAL(Y).EQ.0.0 .AND. AIMAG(Y).EQ.0.0) GO TO 70
    DO 50 K=1,N
      A(J,K) = A(J,K) - Y*A(I,K)
    50  CONTINUE
    DO 60 K=IPL,N
      B(J,K) = B(J,K) - Y*B(I,K)
    60  CONTINUE
  70  CONTINUE
  B(IPL,I) = (0.0,0.0)
  80  CONTINUE
C INITIALIZE X
  IF (.NOT.WANTX) GO TO 110
  DO 100 I=1,N
    DO 90 J=1,N
      X(I,J) = (0.0,0.0)
    90  CONTINUE
    X(I,I) = (1.0,0.00)
  100 CONTINUE
C REDUCE A TO UPPER HESSENBERG FORM
110 NM2 = N - 2
  IF (NM2.LT.1) GO TO 270
  DO 260 J=1,NM2
    JM2 = NM1 - J
    JP1 = J + 1
    DO 250 II=1,JM2
      I = N + 1 - II
      IML = I - 1
      IMJ = I - J
      W = A(I,J)
      Z = A(IML,J)
      IF (ABS(REAL(W))+ABS(AIMAG(W)).LE.ABS(REAL(Z))
        * +ABS(AIMAG(Z))) GO TO 140
C MUST INTERCHANGE ROWS
      DO 120 K=J,N
        Y = A(I,K)
        A(I,K) = A(IML,K)
        A(IML,K) = Y
      120 CONTINUE
      DO 130 K=IML,N
        Y = B(I,K)
        B(I,K) = B(IML,K)
        B(IML,K) = Y
      130 CONTINUE
      140 Z = A(I,J)
      IF (REAL(Z).EQ.0.0 .AND. AIMAG(Z).EQ.0.0) GO TO 170
      Y = Z/A(IML,J)
      DO 150 K=JP1,N
        A(I,K) = A(I,K) - Y*A(IML,K)
      150 CONTINUE
      DO 160 K=IML,N
        B(I,K) = B(I,K) - Y*B(IML,K)
      160 CONTINUE
C TRANSFORMATION FROM THE RIGHT
  170 W = B(I,IML)
      Z = B(I,I)
      IF (ABS(REAL(W))+ABS(AIMAG(W)).LE.ABS(REAL(Z))
        * +ABS(AIMAG(Z))) GO TO 210
C MUST INTERCHANGE COLUMNS
      DO 180 K=1,I
        Y = B(K,I)
        B(K,I) = B(K,IML)
        B(K,IML) = Y
      180 CONTINUE
      DO 190 K=1,N
        Y = A(K,I)
        A(K,I) = A(K,IML)
        A(K,IML) = Y
      190 CONTINUE
      IF (.NOT.WANTX) GO TO 210

```

LZH 310
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LZH 1180
LZH 1190
LZH 1200
LZH 1210
LZH 1220

```

DO 200 K=IMJ,N
  Y = X(K,I)
  X(K,I) = X(K,IM1)
  X(K,IM1) = Y
200 CONTINUE
210 Z = B(I,IM1)
  IF (REAL(Z).EQ.0.0 .AND. AIMAG(Z).EQ.0.0) GO TO 250
  Y = Z/B(I,I)
  DO 220 K=1,IM1
    B(K,IM1) = B(K,IM1) - Y*B(K,I)
220 CONTINUE
  B(I,IM1) = (0.0,0.0)
  DO 230 K=1,N
    A(K,IM1) = A(K,IM1) - Y*A(K,I)
230 CONTINUE
  IF (.NOT.WANTX) GO TO 250
  DO 240 K=IMJ,N
    X(K,IM1) = X(K,IM1) - Y*X(K,I)
240 CONTINUE
250 CONTINUE
  A(JP1+1,J) = (0.0,0.0)
260 CONTINUE
270 RETURN
END
LZH 1230
LZH 1240
LZH 1250
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LZH 1380
LZH 1390
LZH 1400
LZH 1410
LZH 1420
LZH 1430
LZH 1440
LZH 1450
LZH 1460

.SUBROUTINE LZIT(N, A, NA, B, NB, X, NX, WANTX, ITER, EIGA,
* EIGB)
C THIS SUBROUTINE SOLVES THE GENERALIZED EIGENVALUE PROBLEM
C A X = LAMBDA B X
C WHERE A IS A COMPLEX UPPER HESSENBERG MATRIX OF
C ORDER N AND B IS A COMPLEX UPPER TRIANGULAR MATRIX OF ORDER N
C INPUT PARAMETERS
C N ORDER OF A AND B
C A AN N X N UPPER HESSENBERG COMPLEX MATRIX
C NA THE ROW DIMENSION OF THE A MATRIX
C B AN N X N UPPER TRIANGULAR COMPLEX MATRIX
C NB THE ROW DIMENSION OF THE B MATRIX
C X CONTAINS TRANSFORMATIONS TO OBTAIN EIGENVECTORS OF
C ORIGINAL SYSTEM. IF EIGENVECTORS ARE REQUESTED AND QZHESS
C IS NOT CALLED, X SHOULD BE SET TO THE IDENTITY MATRIX
C NX THE ROW DIMENSION OF THE X MATRIX
C WANTX LOGICAL VARIABLE WHICH SHOULD BE SET TO .TRUE.
C IF EIGENVECTORS ARE WANTED. OTHERWISE IT
C SHOULD BE SET TO .FALSE.
C OUTPUT PARAMETERS
C X THE ITH COLUMN CONTAINS THE ITH EIGENVECTOR
C IF EIGENVECTORS ARE REQUESTED
C ITER AN INTEGER ARRAY OF LENGTH N WHOSE ITH ENTRY
C CONTAINS THE NUMBER OF ITERATIONS NEEDED TO FIND
C THE ITH EIGENVALUE. FOR ANY I IF ITER(I) =-1 THEN
C AFTER 30 ITERATIONS THERE HAS NOT BEEN A SUFFICIENT
C DECREASE IN THE LAST SUBDIAGONAL ELEMENT OF A
C TO CONTINUE ITERATING.
C EIGA A COMPLEX ARRAY OF LENGTH N CONTAINING THE DIAGONAL OF A
C EIGB A COMPLEX ARRAY OF LENGTH N CONTAINING THE DIAGONAL OF B
C THE ITH EIGENVALUE CAN BE FOUND BY DIVIDING EIGA(I) BY
C EIGB(I). WATCH OUT FOR EIGB(I) BEING ZERO
C EIGB(I), B(NB,N), EIGA(N), EIGB(N)
C COMPLEX S, W, Y, Z, CSQRT
C COMPLEX X(NX,N)
C INTEGER ITER(N)
C COMPLEX ANNM1, ALFM, BETM, D, SL, DEN, NUM, ANM1M1
C REAL EPSA, EPSB, SS, R, ANORM, BNORM, ANI, BNI, C
C REAL D0, D1, D2, E0, E1
C LOGICAL WANTX
C NN = N
LZI 10
LZI 20
LZI 30
LZI 40
LZI 50
LZI 60
LZI 70
LZI 80
LZI 90
LZI 100
LZI 110
LZI 120
LZI 130
LZI 140
LZI 150
LZI 160
LZI 170
LZI 180
LZI 190
LZI 200
LZI 210
LZI 220
LZI 230
LZI 240
LZI 250
LZI 260
LZI 270
LZI 280
LZI 290
LZI 300
LZI 310
LZI 320
LZI 330
LZI 340
LZI 350
LZI 360
LZI 370
LZI 380
LZI 390
LZI 400
LZI 410
C COMPUTE THE MACHINE PRECISION TIMES THE NORM OF A AND B
ANORM = 0.
BNORM = 0.
DO 30 I=1,N
  ANI = 0.
  IF (I.EQ.1) GO TO 10
  Y = A(I,I-1)
  ANI = ANI + ABS(REAL(Y)) + ABS(AIMAG(Y))
10 BNI = 0.
  DO 20 J=I,N
    ANI = ANI + ABS(REAL(A(I,J))) + ABS(AIMAG(A(I,J)))
    BNI = BNI + ABS(REAL(B(I,J))) + ABS(AIMAG(B(I,J)))
20 CONTINUE
  IF (ANI.GT.ANORM) ANORM = ANI
  IF (BNI.GT.BNORM) BNORM = BNI
30 CONTINUE
  IF (ANORM.EQ.0.) ANORM = 1.0
  IF (BNORM.EQ.0.) BNORM = 1.0
  EPSB = BNORM
  EPSA = ANORM
40 EPSA = EPSA/2.0
  EPSB = EPSB/2.0
  C = ANORM + EPSA
  IF (C.GT.ANORM) GO TO 40
  IF (N.LE.1) GO TO 320
LZI 420
LZI 430
LZI 440
LZI 450
LZI 460
LZI 470
LZI 480
LZI 490
LZI 500
LZI 510
LZI 520
LZI 530
LZI 540
LZI 550
LZI 560
LZI 570
LZI 580
LZI 590
LZI 600
LZI 610
LZI 620
LZI 630
LZI 640
LZI 650
LZI 660

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```

50 ITS = 0
   NML = NN - 1
C CHECK FOR NEGLIGIBLE SUBDIAGONAL ELEMENTS
60 D2 = ABS(REAL(A(NN,NN))) + ABS(AIMAG(A(NN,NN)))
   DO 70 LB=2,NN
     L = NN + 2 - LB
     SS = D2
     Y = A(L-1,L-1)
     D2 = ABS(REAL(Y)) + ABS(AIMAG(Y))
     SS = SS + D2
     Y = A(L,L-1)
     R = SS + ABS(REAL(Y)) + ABS(AIMAG(Y))
     IF (R.EQ.SS) GO TO 80
70 CONTINUE
   L = 1
80 IF (L.EQ.NN) GO TO 320
   IF (ITS.LT.30) GO TO 90
   ITER(NN) = -1
   IF (ABS(REAL(A(NN,NML))) + ABS(AIMAG(A(NN,NML)))) .GT. 0.8 *
     * ABS(REAL(ANNM1)) + ABS(AIMAG(ANNM1))) RETURN
90 IF (ITS.EQ.10 .OR. ITS.EQ.20) GO TO 110
C COMPUTE SHIFT AS EIGENVALUE OF LOWER 2 BY 2
   ANNM1 = A(NN,NML)
   ANM1M1 = A(NML,NML)
   S = A(NN,NN)*B(NML,NML) - ANNM1*B(NML,NN)
   W = ANNM1*B(NN,NN)*(A(NML,NN)*B(NML,NML) - B(NML,NN)*ANM1M1)
   Y = (ANM1M1*B(NN,NN) - S)/2.
   Z = CSQRT(Y*Y+W)
   IF (REAL(Z).EQ.0.0 .AND. AIMAG(Z).EQ.0.0) GO TO 100
   D0 = REAL(Y/Z)
   IF (D0.LT.0.0) Z = -Z
100 DEN = (Y+Z)*B(NML,NML)*B(NN,NN)
   IF (REAL(DEN).EQ.0.0 .AND. AIMAG(DEN).EQ.0.0) DEN =
     * CMLX(EPSA,0.0)
   NUM = (Y+Z)*S - W
   GO TO 120
C AD-HOC SHIFT
110 Y = A(NML,NN-2)
   NUM = CMLX(ABS(REAL(ANNM1)) + ABS(AIMAG(ANNM1)), ABS(REAL(Y))
     * +ABS(AIMAG(Y)))
   DEN = (1.0,0.0)
C CHECK FOR 2 CONSECUTIVE SMALL SUBDIAGONAL ELEMENTS
120 IF (NN.EQ.L+1) GO TO 140
   D2 = ABS(REAL(A(NML,NML))) + ABS(AIMAG(A(NML,NML)))
   E1 = ABS(REAL(ANNM1)) + ABS(AIMAG(ANNM1))
   D1 = ABS(REAL(A(NN,NN))) + ABS(AIMAG(A(NN,NN)))
   NL = NN - (L+1)
   DO 130 MB=1,NL
     M = NN - MB
     E0 = E1
     Y = A(M,M-1)
     E1 = ABS(REAL(Y)) + ABS(AIMAG(Y))
     D0 = D1
     D1 = D2
     Y = A(M-1,M-1)
     D2 = ABS(REAL(Y)) + ABS(AIMAG(Y))
     Y = A(M,M)*DEN - B(M,M)*NUM
     D0 = (D0+D1+D2)*(ABS(REAL(Y)) + ABS(AIMAG(Y)))
     E0 = E0*E1*(ABS(REAL(DEN)) + ABS(AIMAG(DEN))) + D0
     IF (E0.EQ.D0) GO TO 150
130 CONTINUE
140 M = L
150 CONTINUE
   ITS = ITS + 1
   W = A(M,M)*DEN - B(M,M)*NUM
   Z = A(M+1,M)*DEN
   D1 = ABS(REAL(Z)) + ABS(AIMAG(Z))
   D2 = ABS(REAL(W)) + ABS(AIMAG(W))
C FIND L AND M AND SET A=LAM AND B=LBM
   NP1 = N + 1
   LOR1 = L
   NNORN = NN
   IF (.NOT.WANTX) GO TO 160
   LOR1 = 1
   NNOKN = N
160 DO 310 I=M,NM1
   J = I + 1
C FIND ROW TRANSFORMATIONS TO RESTORE A TO
C UPPER HESSENBERG FORM. APPLY TRANSFORMATIONS
C TO A AND B
   IF (I.EQ.M) GO TO 170
   W = A(I,I-1)
   Z = A(J,I-1)
   D1 = ABS(REAL(Z)) + ABS(AIMAG(Z))
   D2 = ABS(REAL(W)) + ABS(AIMAG(W))
   IF (D1.EQ.0.0) GO TO 60
170 IF (D2.GT.D1) GO TO 190
C MUST INTERCHANGE ROWS
   DO 180 K=I,NNORN
     Y = A(I,K)
     A(I,K) = A(J,K)
     A(J,K) = Y
     Y = B(I,K)

```

LZI 670
LZI 680
LZI 690
LZI 700
LZI 710
LZI 720
LZI 730
LZI 740
LZI 750
LZI 760
LZI 770
LZI 780
LZI 790
LZI 800
LZI 810
LZI 820
LZI 830
LZI 840
LZI 850
LZI 860
LZI 870
LZI 880
LZI 890
LZI 900
LZI 910
LZI 920
LZI 930
LZI 940
LZI 950
LZI 960
LZI 970
LZI 980
LZI 990
LZI 1000
LZI 1010
LZI 1020
LZI 1030
LZI 1040
LZI 1050
LZI 1060
LZI 1070
LZI 1080
LZI 1090
LZI 1100
LZI 1110
LZI 1120
LZI 1130
LZI 1140
LZI 1150
LZI 1160
LZI 1170
LZI 1180
LZI 1190
LZI 1200
LZI 1210
LZI 1220
LZI 1230
LZI 1240
LZI 1250
LZI 1260
LZI 1270
LZI 1280
LZI 1290
LZI 1300
LZI 1310
LZI 1320
LZI 1330
LZI 1340
LZI 1350
LZI 1360
LZI 1370
LZI 1380
LZI 1390
LZI 1400
LZI 1410
LZI 1420
LZI 1430
LZI 1440
LZI 1450
LZI 1460
LZI 1470
LZI 1480
LZI 1490
LZI 1500
LZI 1510
LZI 1520
LZI 1530
LZI 1540
LZI 1550
LZI 1560
LZI 1570
LZI 1580
LZI 1590

```

        B(I,K) = B(J,K)
        B(J,K) = Y
180  CONTINUE
        IF (I.GT.M) A(I,I-1) = A(J,I-1)
        IF (D2.EQ.0.0) GO TO 220
C THE SCALING OF W AND Z IS DESIGNED TO AVOID A DIVISION BY ZERO
C WHEN THE DENOMINATOR IS SMALL
        Y = CMPLX(REAL(W)/D1,AIMAG(W)/D1)/CMPLX(REAL(Z)/D1,AIMAG(Z)/
        * D1)
        GO TO 200
190  Y = CMPLX(REAL(Z)/D2,AIMAG(Z)/D2)/CMPLX(REAL(W)/D2,AIMAG(W)/
        * D2)
200  DO 210 K=I,NNORN
        A(J,K) = A(J,K) - Y*A(I,K)
        B(J,K) = B(J,K) - Y*B(I,K)
210  CONTINUE
        IF (I.GT.M) A(J,I-1) = (0.0,0.0)
C PERFORM TRANSFORMATIONS FROM RIGHT TO RESTORE B TO
C TRIANGULAR FORM
C APPLY TRANSFORMATIONS TO A
220  Z = B(J,I)
        W = B(J,J)
        D2 = ABS(REAL(W)) + ABS(AIMAG(W))
        D1 = ABS(REAL(Z)) + ABS(AIMAG(Z))
        IF (D1.EQ.0.0) GO TO 60
        IF (D2.GT.D1) GO TO 270
C MUST INTERCHANGE COLUMNS
        DO 230 K=LOR1,J
        Y = A(K,J)
        A(K,J) = A(K,I)
        A(K,I) = Y
        Y = B(K,J)
        B(K,J) = B(K,I)
        B(K,I) = Y
230  CONTINUE
        IF (I.EQ.NM1) GO TO 240
        Y = A(J+1,J)
        A(J+1,J) = A(J+1,I)
        A(J+1,I) = Y
240  IF (.NOT.WANTX) GO TO 260
        DO 250 K=1,N
        Y = X(K,J)
        X(K,J) = X(K,I)
        X(K,I) = Y
250  CONTINUE
260  IF (D2.EQ.0.0) GO TO 310
        Z = CMPLX(REAL(W)/D1,AIMAG(W)/D1)/CMPLX(REAL(Z)/D1,AIMAG(Z)/
        * D1)
        GO TO 280
270  Z = CMPLX(REAL(Z)/D2,AIMAG(Z)/D2)/CMPLX(REAL(W)/D2,AIMAG(W)/
        * D2)
280  DO 290 K=LOR1,J
        A(K,I) = A(K,I) - Z*A(K,J)
        B(K,I) = B(K,I) - Z*B(K,J)
290  CONTINUE
        B(J,I) = (0.0,0.0)
        IF (I.LT.NM1) A(I+2,I) = A(I+2,I) - Z*A(I+2,J)
        IF (.NOT.WANTX) GO TO 310
        DO 300 K=1,N
        X(K,I) = X(K,I) - Z*X(K,J)
300  CONTINUE
310  CONTINUE
        GO TO 60
320  CONTINUE
        EIGA(NN) = A(NN,NN)
        EIGB(NN) = B(NN,NN)
        IF (NN.EQ.1) GO TO 330
        ITER(NN) = ITS
        NN = NM1
        IF (NN.GT.1) GO TO 50
        ITER(1) = 0
        GO TO 320
C FIND EIGENVECTORS USING B FOR INTERMEDIATE STORAGE
330  IF (.NOT.WANTX) RETURN
        M = N
340  CONTINUE
        ALFM = A(M,M)
        BETM = B(M,M)
        B(M,M) = (1.0,0.0)
        L = M - 1
        IF (L.EQ.0) GO TO 370
350  CONTINUE
        L1 = L + 1
        SL = (0.0,0.0)
        DO 360 J=L1,M
        SL = SL + (BETM*A(L,J) - ALFM*B(L,J))*B(J,M)
360  CONTINUE
        Y = BETM*A(L,L) - ALFM*B(L,L)
        IF (REAL(Y).EQ.0.0 .AND. AIMAG(Y).EQ.0.0) Y =
        * CMPLX((EPSA+EPSB)/2.0,0.0)
        B(L,M) = -SL/Y
        L = L - 1
LZI 1600
LZI 1610
LZI 1620
LZI 1630
LZI 1640
LZI 1650
LZI 1660
LZI 1670
LZI 1680
LZI 1690
LZI 1700
LZI 1710
LZI 1720
LZI 1730
LZI 1740
LZI 1750
LZI 1760
LZI 1770
LZI 1780
LZI 1790
LZI 1800
LZI 1810
LZI 1820
LZI 1830
LZI 1840
LZI 1850
LZI 1860
LZI 1870
LZI 1880
LZI 1890
LZI 1900
LZI 1910
LZI 1920
LZI 1930
LZI 1940
LZI 1950
LZI 1960
LZI 1970
LZI 1980
LZI 1990
LZI 2000
LZI 2010
LZI 2020
LZI 2030
LZI 2040
LZI 2050
LZI 2060
LZI 2070
LZI 2080
LZI 2090
LZI 2100
LZI 2110
LZI 2120
LZI 2130
LZI 2140
LZI 2150
LZI 2160
LZI 2170
LZI 2180
LZI 2190
LZI 2200
LZI 2210
LZI 2220
LZI 2230
LZI 2240
LZI 2250
LZI 2260
LZI 2270
LZI 2280
LZI 2290
LZI 2300
LZI 2310
LZI 2320
LZI 2330
LZI 2340
LZI 2350
LZI 2360
LZI 2370
LZI 2380
LZI 2390
LZI 2400
LZI 2410
LZI 2420
LZI 2430
LZI 2440
LZI 2450
LZI 2460
LZI 2470
LZI 2480
LZI 2490
LZI 2500
LZI 2510

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370 IF (L.GT.0) GO TO 350	LZI 2520
M = M - 1	LZI 2530
IF (M.GT.0) GO TO 340	LZI 2540
C TRANSFORM TO ORIGINAL COORDINATE SYSTEM	LZI 2550
M = N	LZI 2560
380 CONTINUE	LZI 2570
DO 400 I=1,N	LZI 2580
S = (0.0,0.0)	LZI 2590
DO 390 J=1,M	LZI 2600
S = S + X(I,J)*B(J,M)	LZI 2610
390 CONTINUE	LZI 2620
X(I,M) = S	LZI 2630
400 CONTINUE	LZI 2640
M = M - 1	LZI 2650
IF (M.GT.0) GO TO 380	LZI 2660
C NORMALIZE SO THAT LARGEST COMPONENT = 1.	LZI 2670
M = N	LZI 2680
410 CONTINUE	LZI 2690
SS = 0.	LZI 2700
DO 420 I=1,N	LZI 2710
R = ABS(REAL(X(I,M))) + ABS(AIMAG(X(I,M)))	LZI 2720
IF (R.LT.SS) GO TO 420	LZI 2730
SS = R	LZI 2740
D = X(I,M)	LZI 2750
420 CONTINUE	LZI 2760
IF (SS.EQ.0.0) GO TO 440	LZI 2770
DO 430 I=1,N	LZI 2780
X(I,M) = X(I,M)/D	LZI 2790
430 CONTINUE	LZI 2800
440 M = M - 1	LZI 2810
IF (M.GT.0) GO TO 410	LZI 2820
RETURN	LZI 2830
END	LZI 2840

ACM Transactions on Mathematical Software, Vol. 2, No. 4, December 1976, Page 396.

REMARK ON ALGORITHM 496

The LZ Algorithm to Solve the Generalized Eigenvalue Problem for Complex Matrices [F2]

[L.C. Kaufman, *ACM Trans. Math. Software* 1, 3 (Sept. 1975), 271-281]

L.C. Kaufman [Recd 22 April 1976]

Bell Laboratories, Murray Hill, NJ 07974.

The following incorrect lines should be changed:

from

IF(I.GT.M) A(J,I--1)=(0.0,0.0)	LZI 1760
220 Z=B(J,I)	LZI 1800

to

220 IF(I.GT.M) A(J,I--1)=(0.0,0.0)	LZI 1760
Z=B(J,I)	LZI 1800

The following unnecessary line should be deleted:

NP1=N+1	LZI 1360
---------	----------

ALGORITHM 497

Automatic Integration of Functional Differential Equations [D2]

KENNETH W. NEVES

Babcock & Wilcox

Key Words and Phrases: functional differential equations, automatic integration, one-step and multistep methods, time lags and retarded arguments, Volterra integro-differential equations, derivative jump discontinuities

CR Categories: 5.16, 5.17, 5.18

Language: Fortran

DESCRIPTION

This algorithm is a complement to [1] where the theoretical background and development are described.

REFERENCES

1. NEVES, K.W. Automatic integration of functional differential equations: an approach. *ACM Trans. Math. Software* 1, 4 (Dec. 1975), 357-368.

ALGORITHM

	SUBROUTINE DMRODE(F, A, B, HMIN, H, HMAX, N, X0, WK, EPS,	DMR	10
	* JSTART, AA, TT, XX, ZZ, FINIT, ALPHA, WKA, IER, IQ)	DMR	20
C		DMR	30
C F	-AUTOMATIC STEP CHANGE MERSON DIFFERENTIAL EQUATION	DMR	40
C U	- SOLVER MODIFIED TO HANDLE FUNCTIONAL DIFFERENTIAL	DMR	50
C N	- EQUATIONS OF THE FORM	DMR	60
C C	- (1) $DX/DT = F(T, X, Z, I)$	DMR	70
C T	- WHERE F (DESCRIBED BELOW) IS A VECTOR VALUED	DMR	80
C I	- FUNCTION OF DIMENSION N, X AND Z ARE	DMR	90
C O	- N VECTORS SUCH THAT THE ITH COMPONENT OF Z IS	DMR	100
C N	- THE ITH COMPONENT OF X EVALUATED AT ALPHA(X,T,I)	DMR	110
C	- ALPHA IS A VECTOR VALUED FUNCTION (DESCRIBED	DMR	120
C	- BELOW). ALPHA IS A USER SUPPLIED STATE DEPENDENT	DMR	130
C	- RETARDING FUNCTION, AND (1) IS CALLED A	DMR	140
C	- RETARDED ORDINARY DIFF. EQ. (RODE) WITH STATE	DMR	150
C	- DEPENDENT LAG (SDL). NOTE -AS WRITTEN ABOVE	DMR	160
C	- EACH COMPONENT OF DX/DT MAY DEPEND ON A DIFFERENT	DMR	170
C	- RETARDING FUNCTION. IT IS ALSO POSSIBLE TO ALLOW	DMR	180
C	- A GIVEN COMPONENT TO DEPEND ON MORE THAN ONE	DMR	190
C	- RETARDING FUNCTION BY AUGMENTING THE ORDER OF	DMR	200
C	- THE SYSTEM (SEE DISCRPTION FOR DETAILS).	DMR	210
C P	F-USER SUPPLIED EXTERNAL FUNCTION. THE DERIVATIVE	DMR	220
C A	- OF THE ITH COMPONENT OF X IS $F(T, X, Z, I)$, WHERE	DMR	230
C R	- Z AND X ARE ARRAYS OF DIMENSION N. X(I)	DMR	240

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C A - IS THE SOLUTION OF THE ITH COMPONENT AT T AND DMR 250
C M - Z (I) IS THE ITH COMPONENT OF X AT ALPHA(X,T,I) DMR 260
C E - Z(I) IS CALCULATED INTERNALLY WITHIN DMRODE AND DMR 270
C T - PASSED TO F (SEE ALPHA BELOW). T IS THE CURRENT DMR 280
C E - VALUE OF THE INDEPENDENT VARIABLE. DMR 290
C R A-STARTING POINT OF INTEGRATION DMR 300
C S B-FINAL POINT OF INTEGRATION (NOT NEC. .GT. A) DMR 310
C HMIN-THE ABSOLUTE VALUE OF MINIMUM ALLOWABLE STEP SIZE DMR 320
C - THE STEP SIZE WILL NEVER BE SMALLER THAN HMIN DMR 330
C - IN ABSOLUTE VALUE EXCEPT AT THE END OF THE DMR 340
C - INTERVAL WHERE H IS ALWAYS .GT. HMIN/2 IN ABS.. DMR 350
C H-INITIAL GUESS OF THE INTEGRATION STEP SIZE. THE DMR 360
C - ADJUSTED STEPSIZE IS RETURNED IN H. DMR 370
C HMAX-THE ABSOLUTE VALUE OF MAXIMUM ALLOWABLE STEP SIZE DMR 380
C N-NUMBER OF SIMULTANEOUS EQUATIONS DMR 390
C X0-VECTOR OF LENGTH N CONTAINING THE VALUE OF X(A). DMR 400
C - X(B) IS RETURNED IN X0. DMR 410
C WK-WORK AREA OF DIMENSION .GE. 4*N DMR 420
C EPS-INPUT REQUEST OF RELATIVE ERROR IN LARGEST COMP. DMR 430
C JSTART-A PARAMETER USED TO INITIALIZE Z ARRAY INTERNALLY DMR 440
C - IN DMRODE. JSTART MUST BE SET EQUAL TO ZERO ON DMR 450
C - FIRST CALL TO DMRODE. TO INTEGRATE BEYOND B AFTER DMR 460
C - FIRST CALL CERTAIN ARRAYS MUST BE SAVED. THIS DMR 470
C - IS DONE AUTOMATICALLY BY LETTING JSTART=1. DMR 480
C - IF JSTART=-1, THEN THE INITIAL CONDITIONS AT DMR 490
C - THE PREVIOUS CALL ARE REINSTATED AND THE INTE- DMR 500
C - GRATION FROM A TO POSSIBLY A NEW CHOICE FOR B DMR 510
C - CAN BE EXECUTED. THIS ALLOWS INTERACTIVE USE OF DMR 520
C - DMRODE IN ORDER TO CALCULATE DERIVATIVE JUMP DMR 530
C - POINTS IN THE SOLUTION(SEE DESCRIPTION) DMR 540
C AA-A PARAMETER THAT MUST BE SET EQUAL TO THE VALUE OF DMR 550
C - THE STARTING POINT A WHEN FIRST CALL TO DMRODE DMR 560
C - WAS MADE (I.E. WHEN JSTART=0) (SEE FINIT) DMR 570
C TT-STORAGE ARRAY FOR SAVING TIME STEP VALUES DMR 580
C - TT IS OF DIMENSION IQ WHERE IQ IS .GT.LENGTH OF DMR 590
C - ENTIRE INTERVAL/HMIN DMR 600
C XX,ZZ-STORAGE ARRAYS OF DIMENSION N BY IQ USED FOR DMR 610
C - SAVING VALUES OF X AND Z DESCRIBED IN F. DMR 620
C FINIT-USER SUPPLIED EXTERNAL FUNCTION, FINIT(T,I). FINIT DMR 630
C - RETURNS THE VALUE OF THE ITH COMPONENT OF DMR 640
C - THE INITIAL FUNCTION AT T WHENEVER ((T.LE.AA) DMR 650
C - .AND.(H.GE.0)).OR.((T.GE.AA).AND.(H.LE.0)). DMR 660
C ALPHA-USER SUPPLIED EXTERNAL FUNCTION OF THE DMR 670
C - FORM ALPHA(X,T,I) WHERE X IS ASSUMED TO BE DMR 680
C - THE CURRENT SOLUTION (SUPPLIED BY DMRODE) AND DMR 690
C - T THE CURRENT TIME. IN GENERAL ALPHA WILL BE DMR 700
C - DIFFERENT FOR EACH COMPONENT X(I). IN THE EVENT DMR 710
C - A RODE IS POSED WITH MORE THAN ONE LAG IN ANY DMR 720
C - GIVEN COMPONENT THE SYSTEM MAY BE AUGMENTED TO DMR 730
C - HANDLE THIS SITUATION (SEE DESCRIPTION). DMR 740
C WKA-WORK AREA OF DIMENSION .GE.11*N UNLESS DMR 750
C - IT IS KNOWN THE LAG T-ALPHA(T,X(T),I) IN ALL DMR 760
C - COMPONENTS WILL ALWAYS EXCEED HMAX OR VANISH DMR 770
C - OVER THE ENTIRE INTERVAL IN A GIVEN COMPONENT, DMR 780
C - IN WHICH CASE A DIMENSION OF 3*N WILL SUFFICE DMR 790
C IER-OUTPUT PARAMETER IF IER=0 AND HMAX.NE.HMIN DMR 800
C - NORMAL ERROR CHECKING TOOK PLACE. IF IER=1 DMR 810
C - HMIN WAS ATTAINED BY THE STEP CHANGING DMR 820
C - PROCEDURE, HENCE REQUESTED ACCURACY IS NOT DMR 830
C - GUARANTEED DMR 840
C IQ-THE DIMENSION OF TT IN THE CALLING PROGRAM (ALSO DMR 850
C - THE DIMENSION OF THE 2ND ARGUMENT OF THE XX AND DMR 860
C - ZZ ARRAYS). DMR 870
C PRECISION - SINGLE DMR 880
C AUTHOR/IMPLEMENTOR- KENNETH W. NEVES DMR 890
C REQUIRED ROUTINES DMR 900
C - DZETA, AN INTERPOLATION SCHEME OF ORDER COMP- DMR 910
C - ARABLE TO 4TH ORDER MERSON METHOD. DZETA IN TURN DMR 920
C - WILL CALL STARTER METHOD WHEN LAG IS SMALLER DMR 930
C - THAN CURRENT STEP SIZE. DMR 940
C DMR 950
C DIMENSION WK(1), X0(1) DMR 960
C DIMENSION WKA(1) DMR 970
C DIMENSION TT(1), XX(N,1), ZZ(N,1) DMR 980
C EXTERNAL F, FINIT, ALPHA DMR 990
C INTEGER Sw DMR 1000
C LOGICAL BE, BH, BR, BX, BT DMR 1010
C DATA ZERO, P5, OP5, THREE, FOUR /0.,.5,1.5,3.,4./ DMR 1020
C E5 = .5*EPS DMR 1030
C IER = 0 DMR 1040
C IF (A.EQ.B) GO TO 290 DMR 1050
C IB1 = N + N DMR 1060
C IB2 = IB1 + N DMR 1070
C ***** DMR 1080
C CHECK FOR THE PROPER SIGN OF H DMR 1090
C ***** DMR 1100
C H = SIGN(ABS(H),B-A) DMR 1110
C ***** DMR 1120
C CHECK FOR 1ST CALL TO DMRODE. IF DMR 1130
C YES, THEN INITIALIZE NEC. ARRAYS DMR 1140
C ***** DMR 1150

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IF (JSTART.NE.0) GO TO 20
JRAY = 1
JSAVE = JRAY
TT(1) = A
DO 10 J=1,N
  XX(J,1) = X0(J)
  ZZ(J,1)=DZETA(TT(1),X0,TT,XX,ZZ,AA,F,FINIT,ALPHA,
* N,J,WKA(N+1),H,1,IQ)
10 CONTINUE
20 CONTINUE
  BH = .TRUE.
  BR = .TRUE.
  BX = .TRUE.
  BT = .TRUE.
  X = A
  XS = A
  IF (ABS(H).GE.HMAX) H = SIGN(1.,H)*HMAX
  DO 30 J=1,N
    IJK0 = N + J
    WK(IJK0) = X0(J)
30 CONTINUE
C *****
C CHECK JSTART AND EXECUTE
C PROPER INITIALIZATION
C *****
IF (JSTART) 40, 90, 50
40 JRAY = JSAVE
50 DO 60 J=1,N
  IJK0 = N + J
  WK(IJK0) = XX(J,JRAY)
60 CONTINUE
  JSAVE = JRAY
  GO TO 90
70 XS = X
C *****
C AFTER EACH SUCCESSFUL TIME STEP
C UPDATE TT,XX,ZZ ARRAYS. CALL TO
C DZETA PERFORMS NECESSARY INTERPO-
C LATION AT LAG POINT.
C *****
TT(JRAY+1) = X
DO 80 J=1,N
  IJK0 = N + J
  WK(IJK0) = X0(J)
  XX(J,JRAY+1) = X0(J)
  HH = X - TT(JRAY)
  ZZ(J,JRAY+1) = DZETA(TT(JRAY+1),X0,TT,XX,ZZ,AA,F,FINIT,
* ALPHA,N,J,WKA(N+1),HH,JRAY,IQ)
80 CONTINUE
C *****
C INSERT CALL TO PRINT ROUTINE HERE
C IF DESIRED.
C *****
JRAY = JRAY + 1
IF (BR) GO TO 90
H = HSS
RETURN
90 HSS = H
C *****
C TEST FOR END OF PROBLEM
C *****
Q = X + H - B
BE = .TRUE.
IF (.NOT.((H.GT.ZERO .AND. Q/H.GE.-.5) .OR. (H.LT.ZERO .AND.
* Q/(-H).LE..5))) GO TO 110
IF (BT) HSS = H
H = .5*(B-X)
IF (.NOT.BT) H = B - X
IF (BT) GO TO 100
BR = .FALSE.
100 BT = .FALSE.
110 CONTINUE
H3 = H/THREE
C *****
C CALCULATE SOLN. AT X+H
C *****
DO 270 SW=1,5
C *****
C DZETA CALCULATES LAG AND SOLUTION
C AT PAST POINT (NEC. FOR F)
C *****
DO 120 JS=1,N
  WKA(JS) = DZETA(X,X0,TT,XX,ZZ,AA,F,FINIT,ALPHA,N,JS,
* WKA(N+1),H,JRAY,IQ)
120 CONTINUE
DO 130 JS=1,N
  WK(JS) = F(X,X0,WKA,JS)
130 CONTINUE
DO 230 I=1,N
  Q = H3*WK(I)
  IJK0 = N + I

```

DMR 1160
DMR 1170
DMR 1180
DMR 1190
DMR 1200
DMR 1210
DMR 1220
DMR 1230
DMR 1240
DMR 1250
DMR 1260
DMR 1270
DMR 1280
DMR 1290
DMR 1300
DMR 1310
DMR 1320
DMR 1330
DMR 1340
DMR 1350
DMR 1360
DMR 1370
DMR 1380
DMR 1390
DMR 1400
DMR 1410
DMR 1420
DMR 1430
DMR 1440
DMR 1450
DMR 1460
DMR 1470
DMR 1480
DMR 1490
DMR 1500
DMR 1510
DMR 1520
DMR 1530
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DMR 1580
DMR 1590
DMR 1600
DMR 1610
DMR 1620
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DMR 1670
DMR 1680
DMR 1690
DMR 1700
DMR 1710
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DMR 1860
DMR 1870
DMR 1880
DMR 1890
DMR 1900
DMR 1910
DMR 1920
DMR 1930
DMR 1940
DMR 1950
DMR 1960
DMR 1970
DMR 1980
DMR 1990
DMR 2000
DMR 2010
DMR 2020
DMR 2030
DMR 2040
DMR 2050
DMR 2060

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      IJK1 = IB1 + I
      IJK2 = IB2 + I
      GO TO (140, 150, 160, 170, 180), SW
140  R = Q
      WK(IJK1) = Q
      GO TO 190
C
150  R = P5*(Q+WK(IJK1))
      GO TO 190
C
160  R = THREE*Q
      WK(IJK2) = R
      R = .375*(R+WK(IJK1))
      GO TO 190
C
170  R = WK(IJK1) + FOUR*Q
      WK(IJK1) = R
      R = OP5*(R-WK(IJK2))
      GO TO 190
C
180  R = P5*(Q+WK(IJK1))
      Q = ABS(R+R-OP5*(Q+WK(IJK2)))
190  X0(I) = WK(IJK0) + R
      IF (SW.NE.5) GO TO 230
      *****
      AUTOMATIC STEP CHANGE
      *****
C
      E = ABS(X0(I))
      R = E5
      IF (E.GE.1.E-3) R = E*E5
      IF (HMIN.EQ.HMAX) GO TO 280
      *****
      TEST ADJUSTMENT OF THE STEP
      *****
C
      IF (Q.LT.R .OR. (.NOT.BX)) GO TO 220
      BR = .TRUE.
      BT = .TRUE.
      BH = .FALSE.
      H = P5*H
      IF (ABS(H).GT.HMIN) GO TO 200
      IF (.NOT.BR) GO TO 280
      *****
      THE STEP IS HALVED RESTORE X AND X0,
      AND GO BACK FOR REPEATED INTEGRATION
      WITH THIS NEW STEP
      *****
C
      H = SIGN(1.,H)*HMIN
      BX = .FALSE.
      IER = 1
200  DO 210 J=1,N
      IJK0 = N + J
      X0(J) = WK(IJK0)
210  CONTINUE
      X = XS
      GO TO 90
C
220  IF (Q.GE.0.03125*R) BE = .FALSE.
230  CONTINUE
      GO TO (240, 270, 250, 260, 270), SW
240  X = X + H3
      GO TO 270
C
250  X = X + P5*H3
      GO TO 270
C
260  X = X + P5*H
270  CONTINUE
      *****
      TEST A POSSIBLE DOUBLING OF THE STEP
      *****
C
      IF (.NOT.(BE .AND. BH .AND. BR)) GO TO 280
      H = H + H
      BX = .TRUE.
280  BH = .TRUE.
      GO TO 70
C
290  DO 300 I=1,N
      X0(I) = ZERO
300  CONTINUE
      RETURN
      END
      *****
      FUNCTION DZETA(TA, XA, T, X, XL, TAU, F, FINIT, ALPHA, N,
      * JCOMP, W, HP, JRAY, NN)
C
C F - INTERPOLATION ROUTINE. CALLED BY DMRODE AND
C U - REQUIRES NO USER INTERFACE OTHER THAN USER
C N - SUPPLIED EXTERNAL FUNCTION ALPHA (SEE DMRODE).
C C - DZETA CALCULATES THE LAG(OR RETARDATION) VIA

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DMR 2070
DMR 2080
DMR 2090
DMR 2100
DMR 2110
DMR 2120
DMR 2130
DMR 2140
DMR 2150
DMR 2160
DMR 2170
DMR 2180
DMR 2190
DMR 2200
DMR 2210
DMR 2220
DMR 2230
DMR 2240
DMR 2250
DMR 2260
DMR 2270
DMR 2280
DMR 2290
DMR 2300
DMR 2310
DMR 2320
DMR 2330
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DMR 2350
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DMR 2370
DMR 2380
DMR 2390
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DMR 2650
DMR 2660
DMR 2670
DMR 2680
DMR 2690
DMR 2700
DMR 2710
DMR 2720
DMR 2730
DMR 2740
DMR 2750
DMR 2760
DMR 2770
DMR 2780
DMR 2790
DMR 2800
DMR 2810
DMR 2820
DMR 2830
DMR 2840
DMR 2850
DMR 2860
DMR 2870

DZE 10
DZE 20
DZE 30
DZE 40
DZE 50
DZE 60
DZE 70

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C T - ALPHA WITH TA, XA AS INPUT, THEN USES 2-POINT DZE 80
C I - HERMITE INTERPOLATION TO APPROXIMATE SOLUTION DZE 90
C O - AT LAG. IF STEP SIZE IS LARGER THAN LAG, HERMITE DZE 100
C N - INTERPOLANT CANNOT BE OBTAINED, SO DZETA CALLS DZE 110
C - A TRUE ONESTEP FDE SOLVER. SINCE DMRODE CALLS DZE 120
C - DZETA THE PARAMETER DESCRIPTION CAN BE FOUND DZE 130
C - IN DMRODE DZE 140
C DIMENSION W(1), XA(1) DZE 150
  DIMENSION T(1), X(N,1), XL(N,1) DZE 160
  EXTERNAL F, FINIT, ALPHA DZE 170
  DATA J /1/ DZE 180
  IF ((J.GE.JRAY) .AND. (J.NE.1)) J = JRAY - 1 DZE 190
  IF (J.LE.0) J = 1 DZE 200
C ***** DZE 210
C CALCULATE LAG DZE 220
C ***** DZE 230
C B = ALPHA(XA, TA, JCOMP) DZE 240
  IF (B.EQ.TA) DZETA = XA(JCOMP) DZE 250
  IF (B.EQ.TA) RETURN DZE 260
  IF ((HP.LT.0.) .AND. (B.GE.TAU)) .OR. ((HP.GE.0.) .AND. DZE 270
* (B.LE.TAU)) GO TO 100 DZE 280
C ***** DZE 290
C SEARCH T ARRAY IN ORDER TO FIND DZE 300
C INTERVAL CONTAINING ALPHA(XA, TA, I).. DZE 310
C IF NO SUCH INTERVAL IS FOUND CALL DZE 320
C START OR IF ARRAY EXCEEDS DZE 330
C DIMENSION IQ OF DMRODE, STOP DZE 340
C ***** DZE 350
  DX = B - T(J) DZE 360
  IF (HP.LT.0.) DX = -DX DZE 370
  IF (DX) 10, 40, 30 DZE 380
10 IF (J.EQ.1) GO TO 40 DZE 390
  J = J - 1 DZE 400
  DX = B - T(J) DZE 410
  IF (HP.LT.0.) DX = -DX DZE 420
  IF (DX) 10, 40, 40 DZE 430
20 J = J + 1 DZE 440
30 IF (J.EQ.NN) GO TO 50 DZE 450
  IF (J+1.GT.JRAY) GO TO 90 DZE 460
  DDX = B - T(J+1) DZE 470
  IF (HP.LT.0.) DDX = -DDX DZE 480
  IF (DDX) 40, 20, 20 DZE 490
40 IF (J.EQ.NN) GO TO 50 DZE 500
  GO TO 60 DZE 510
C DZE 520
50 STOP DZE 530
60 CONTINUE DZE 540
C ***** DZE 550
C 2-POINT HERMITE INTERPOLATION DZE 560
C ***** DZE 570
  IF (T(J+1).GE.TA) GO TO 90 DZE 580
  DX = T(J+1) - T(J) DZE 590
  HH = B - T(J) DZE 600
  DO 70 I=1,N DZE 610
    II = I + N DZE 620
    W(I) = X(I,J) DZE 630
    W(II) = XL(I,J) DZE 640
70 CONTINUE DZE 650
  AC = F(T(J), W, W(N+1), JCOMP) DZE 660
  DO 80 I=1,N DZE 670
    II = I + N DZE 680
    W(I) = X(I, J+1) DZE 690
    W(II) = XL(I, J+1) DZE 700
80 CONTINUE DZE 710
  AF2P = F(T(J+1), W, W(N+1), JCOMP) DZE 720
  DIVDF1 = (X(JCOMP, J+1) - X(JCOMP, J)) / DX DZE 730
  DIVDF3 = AC + AF2P - 2.*DIVDF1 DZE 740
  C3 = (DIVDF1 - AC - DIVDF3) / DX DZE 750
  C4 = DIVDF3 / DX**2 DZE 760
  DZETA = X(JCOMP, J) + HH*(AC + HH*(C3 + HH*C4)) DZE 770
  GO TO 110 DZE 780
C DZE 790
90 CONTINUE DZE 800
  J = JRAY DZE 810
  CALL START(TA, XA, T, X, XL, F, FINIT, ALPHA, N, JCOMP, W, J, DZE 820
* ANS, HP) DZE 830
  DZETA = ANS DZE 840
  RETURN DZE 850
C DZE 860
100 DZETA = FINIT(B, JCOMP) DZE 870
110 CONTINUE DZE 880
  RETURN DZE 890
C DZE 900
C DZE 910
  END DZE 920

SUBROUTINE START(TA, XA, T, X, XL, F, FINIT, ALPHA, N, JCOMP, STA 10
* W, JEND, ANS, H1) STA 20
C F -START IS A SUBROUTINE THAT CAN APPROXIMATE THE STA 30

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C U - SOLUTION AT A PAST POINT GIVEN ONLY THE SOLUTION STA 40
C N - AT THE MOST RECENT MESH POINT AND THE FUNCTIONS STA 50
C C - F AND ALPHA, AS DESCRIBED BY DMRODE. IT IS A STA 60
C T - RUNGE-KUTTA TYPE ALGORITHM MODIFIED FOR RODES STA 70
C I - AND IS GLOBALLY 3RD ORDER, BUT PREDICTS 4TH STA 80
C O - ORDER APPROX. TO BE USED IN EVALUATION OF F STA 90
C N - IN MERSON METHOD OF DMRODE. START, DZETA, AND STA 100
C - DMRODE COMBINE TO GIVE 4TH ORDER SIMPLIFIED STA 110
C - PREDICTOR-CORRECTOR TWO STEP ALGORITHM, WHERE STA 120
C - START IS A TRUE ONE-STEP PREDICTOR USED ONLY STA 130
C - WHEN STEP SIZE EXCEEDS THE LAG STA 140
DIMENSION T(1), X(N,1), XL(N,1), XA(1), W(1) STA 150
I1 = N STA 160
I2 = I1 + N STA 170
I3 = I2 + N STA 180
I4 = I3 + N STA 190
I5 = I4 + N STA 200
I6 = I5 + N STA 210
I7 = I6 + N STA 220
I8 = I7 + N STA 230
I9 = I8 + N STA 240
JRAY = JEND STA 250
C ***** STA 260
C RECALCULATE LAG STA 270
C ***** STA 280
B = ALPHA(XA,TA,JCOMP) STA 290
10 TEST = B - T(JRAY) STA 300
IF (H1.LT.0.) TEST = -TEST STA 310
IF (TEST.GT.0.) GO TO 20 STA 320
JRAY = JRAY - 1 STA 330
GO TO 10 STA 340
C ***** STA 350
C CONSTRUCT EULER-TYPE PREDICTOR STA 360
C CALCULATE APPROX. SOLN. AT TA+H STA 370
C SEARCH FOR LAG CORRESPONDING TO STA 380
C X(TA+H) JUST CALCULATED. STA 390
C APPROXIMATE X AT LAG VIA EULER STA 400
C ***** STA 410
20 CONTINUE STA 420
H = H1 STA 430
ISWITC = 1 STA 440
IF (JRAY.NE.JEND) H = T(JRAY+1) - T(JRAY) STA 450
DO 30 I=1,N STA 460
L5 = I5 + I STA 470
L6 = I6 + I STA 480
W(L5) = X(I,JRAY) STA 490
W(L6) = XL(I,JRAY) STA 500
30 CONTINUE STA 510
DO 40 I=1,N STA 520
L1 = I1 + I STA 530
W(L1) = F(T(JRAY),W(I5+1),W(I6+1),I) STA 540
40 CONTINUE STA 550
DO 50 I=1,N STA 560
L5 = I5 + I STA 570
L1 = I1 + I STA 580
W(L5) = X(I,JRAY) + H*W(L1) STA 590
C ***** STA 600
C IN EACH COMPONENT CALCULATE THE STA 610
C CORRESPONDING LAG,W(I), FOR THE STA 620
C EULER PREDICTED VALUE W(L5), AND STA 630
C EXECUTE SEARCH. STA 640
C ***** STA 650
50 CONTINUE STA 660
DO 90 I=1,N STA 670
W(I) = ALPHA(W(I5+1),T(JRAY)+H,I) STA 680
J = JRAY STA 690
60 TEST = W(I) - T(J) STA 700
IF (H1.LT.0.) TEST = -TEST STA 710
IF (TEST.GE.0.) GO TO 70 STA 720
J = J - 1 STA 730
IF (J.EQ.0) GO TO 70 STA 740
GO TO 60 STA 750
C STA 760
70 CONTINUE STA 770
L9 = I9 + I STA 780
W(L9) = J STA 790
IF (J.EQ.0) GO TO 90 STA 800
DO 80 II=1,N STA 810
L3 = I3 + II STA 820
L4 = I4 + II STA 830
W(L3) = X(II,J) STA 840
W(L4) = XL(II,J) STA 850
80 CONTINUE STA 860
90 CONTINUE STA 870
DO 100 I=1,N STA 880
L9 = I9 + I STA 890
J = W(L9) STA 900
H = H1 STA 910
L6 = I6 + I STA 920
IF (J.EQ.0) W(L6) = FINT(W(I),I) STA 930

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          IF (J.EQ.0) GO TO 100
          IF (J.NE.JEND) H = T(J+1) - T(J)
C          *****
C          W(L6) IS THE SOLUTION AT ALPHA(W(I),T,I)
C          *****
          W(L6) = X(I,J) + (W(I)-T(J))*F(T(J),W(I3+1),W(I4+1),I)
100 CONTINUE
          *****
          REPEAT ENTIRE PROCEDURE ABOVE
          USING EULER VALUE TO CALCULATE
          HEUN-TYPE 2ND ORDER PREDICTOR
          *****
          DO 110 I=1,N
            L2 = I2 + I
            W(L2) = F(T(JRAY)+H,W(I5+1),W(I6+1),I)
110 CONTINUE
          DO 120 I=1,N
            L1 = I1 + I
            L2 = I2 + I
            L3 = I3 + I
            L7 = I7 + I
          *****
          FOR EACH COMPONENT ESTIMATE THE
          SOLUTION AT T+H AND T+H/2 AND STORE
          ANSWERS IN W(L3) AND W(L7) RESP.
          *****
          W(L3) = X(I,JRAY) + H*W(L1) + .5*H*(W(L2)-W(L1))
          W(L7) = X(I,JRAY) + .5*H*W(L1) + .125*H*(W(L2)-W(L1))
          *****
          IN STATEMENTS 120 THRU 280 THE COR-
          RESPONDING LAGS AND THE SOLUTION AT
          THESE LAGS ARE COMPUTED TO 2ND ORDER
          BY HEUN METHOD AND STORED IN W(L4) AND
          W(L8) RESPECTIVELY
          *****
120 CONTINUE
          DO 130 I=1,N
            W(I) = ALPHA(W(I3+1),T(JRAY)+H,I)
130 CONTINUE
140 DO 170 I=1,N
          J = JRAY
150 TEST = W(I) - T(J)
          IF (H1.LT.0.) TEST = -TEST
          IF (TEST.GE.0.) GO TO 160
          J = J - 1
          IF (J.EQ.0) GO TO 160
          GO TO 150
C
160 CONTINUE
          L9 = I9 + I
          W(L9) = J
170 CONTINUE
          DO 290 I=1,N
            L9 = I9 + I
            J = W(L9)
            H = H1
            IF (J.EQ.0) GO TO 250
            IF (J.NE.JEND) H = T(J+1) - T(J)
            IF (J.EQ.JEND) GO TO 220
            DO 180 II=1,N
              L5 = I5 + II
              L6 = I6 + II
              W(L5) = X(II,J)
              W(L6) = XL(II,J)
180 CONTINUE
            DO 190 II=1,N
              L1 = I1 + II
              W(L1) = F(T(J),W(I5+1),W(I6+1),II)
190 CONTINUE
            DO 200 II=1,N
              L5 = I5 + II
              L6 = I6 + II
              W(L5) = X(II,J+1)
              W(L6) = XL(II,J+1)
200 CONTINUE
            DO 210 II=1,N
              L2 = I2 + II
              W(L2) = F(T(J+1),W(I5+1),W(I6+1),II)
210 CONTINUE
220 GO TO (230, 240), ISWITC
230 L1 = I1 + I
          L2 = I2 + I
          L4 = I4 + I
          W(L4) = X(I,J) + (W(I)-T(J))*W(L1) + .5*(W(I)-T(J))**2*
          * (W(L2)-W(L1))/H
          GO TO 260
C
240 L1 = I1 + I
          L2 = I2 + I
          L8 = I8 + I

```

```

STA 940
STA 950
STA 960
STA 970
STA 980
STA 990
STA 1000
STA 1010
STA 1020
STA 1030
STA 1040
STA 1050
STA 1060
STA 1070
STA 1080
STA 1090
STA 1100
STA 1110
STA 1120
STA 1130
STA 1140
STA 1150
STA 1160
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STA 1770
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STA 1790
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STA 1810
STA 1820
STA 1830

```

```

      W(L8) = X(I,J) + (W(I)-T(J))*W(L1) + .5*(W(I)-T(J))**2*
      * (W(L2)-W(L1))/H
250  CONTINUE
      I4 = I4 + I
      IF (ISWITC.EQ.1) W(L4) = FINT(W(I),I)
      L8 = I8 + I
      IF (ISWITC.EQ.2) W(L8) = FINT(W(I),I)
      GO TO 260
C
260  GO TO (270, 290), ISWITC
270  DO 280 II=1,N
      W(II) = ALPHA(W(I7+1),T(JRAY)+.5*H,II)
280  CONTINUE
      ISWITC = ISWITC + 1
      GO TO 140
C
C
C
C
C
      *****
      USE HEUN PREDICTOR FOR 4TH ORDER
      APPROX. SOLUTION.
      *****
290  CONTINUE
      H = H1
      IF (JRAY.NE.JEND) H = T(JRAY+1) - T(JRAY)
      D = F(T(JRAY)+H,W(I3+1),W(I4+1),JCOMP)
      C = F(T(JRAY)+.5*H,W(I7+1),W(I8+1),JCOMP)
      SS = B - T(JRAY)
      R = SS/H
      DO 300 I=1,N
      L5 = I5 + I
      L6 = I6 + I
      W(L5) = X(I,JRAY)
      W(L6) = XL(I,JRAY)
300  CONTINUE
      I = JCOMP
      L1 = I1 + I
      W(L1) = F(T(JRAY),W(I5+1),W(I6+1),JCOMP)
C
C
C
C
      *****
      RETURN 4TH ORDER APROXIMATION FOR
      USE IN F IN MERSON METHOD
      *****
      ANS = X(I,JRAY) + SS*(W(L1)+.5*R*(-3.*W(L1)+4.*C-D)+R*R/3.*
      * (2.*W(L1)-4.*C+2.*D))
      RETURN
C
      END

```

```

STA 1840
STA 1850
STA 1860
STA 1870
STA 1880
STA 1890
STA 1900
STA 1910
STA 1920
STA 1930
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STA 1950
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STA 1970
STA 1980
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STA 2010
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STA 2100
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```

ALGORITHM 498

Airy Functions Using Chebyshev Series Approximations

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Key Words and Phrases: Airy functions, Chebyshev series approximation, Chebyshev coefficients, asymptotic expansion, Taylor expansion

CR Categories: 5.12

Language: Fortran

DESCRIPTION

Purpose

This subroutine evaluates the Airy functions $Ai(z)$ and $Bi(z)$ and their respective derivatives $Ai'(z)$ and $Bi'(z)$, for all real values of z within computer capability, by means of Chebyshev series approximations.

Method

(a) For $z \leq -7$ the following asymptotic expansions are used (see [1, eqs. 10.4.60, 10.4.64, 10.4.62, and 10.4.67]):

$$Ai(-z) \sim z^{-1/4}\{sf(z) - cq(z)\}, \quad Ai'(-z) \sim -z^{1/4}\{cp(z) + sq(z)\}$$

$$Bi(-z) \sim z^{-1/4}\{cf(z) + sg(z)\}, \quad Bi'(-z) \sim z^{1/4}\{sp(z) - cq(z)\}$$

where $s = \sin(\zeta + \pi/4)$, $c = \cos(\zeta + \pi/4)$, $\zeta = 2/3 z^{3/2}$, and f , g , p , and q are approximated by the following Chebyshev expansions:

$$f \simeq \sum_{r=0}^{m_1} a_r T_r^*(t), \quad g \simeq \zeta^{-1} \sum_{r=0}^{m_2} b_r T_r^*(t), \quad p \simeq \sum_{r=0}^{m_3} c_r T_r^*(t), \quad q \simeq \zeta^{-1} \sum_{r=0}^{m_4} d_r T_r^*(t)$$

where $t = -(7/z)^{2/3}$, the T_r^* are the shifted Chebyshev polynomials, and the a_r , b_r , c_r , and d_r are the prescribed Chebyshev coefficients.

(b) For $-7 < z \leq 0$ the functions are given by the following Taylor expansions (see [1, eqs. 10.4.2 and 10.4.3]):

$$Ai = e_1 f(z) - e_2 g(z), \quad Ai' = e_1 f'(z) - e_2 g'(z)$$

$$Bi = \sqrt{3}\{e_1 f(z) + e_2 g(z)\}, \quad Bi' = \sqrt{3}\{e_1 f'(z) + e_2 g'(z)\}$$

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where e_1 and e_2 are prescribed constants, and

$$f \simeq \sum_{r=0}^{m_1} a_r T_r^*(t), \quad g \simeq z \sum_{r=0}^{m_2} b_r T_r^*(t)$$

$$f' \simeq z^2 \sum_{r=0}^{m_3} c_r T_r^*(t), \quad g' \simeq \sum_{r=0}^{m_4} d_r T_r^*(t)$$

where $t = -(z/7)^3$.

(c) For $0 < z < 7$ the previous Taylor expansions do not allow the functions to be accurately calculated (unless using very high precision) at all points of the range because of the large fluctuation in magnitude of the functions over this range. To overcome this, the following Chebyshev approximations to the weighted functions are used:

$$\text{Ai exp}(1.75z) \simeq \sum_{r=0}^{m_1} a_r T_r^*(t), \quad \text{Ai}' \text{ exp}(1.75z) \simeq \sum_{r=0}^{m_3} c_r T_r^*(t)$$

$$\text{Bi exp}(-1.75z) \simeq \sum_{r=0}^{m_2} b_r T_r^*(t), \quad \text{Bi}' \text{ exp}(-1.75z) \simeq \sum_{r=0}^{m_4} d_r T_r^*(t)$$

where $t = z/7$.

(d) For $z \geq 7$ the following asymptotic approximations are used (see [1, eqs. 10.4.59, 10.4.63, 10.4.61, and 10.4.66]):

$$\text{Ai}(z) \sim z^{-1/4} \exp(-\zeta) \sum_{r=0}^{m_1} a_r T_r^*(t), \quad \text{Ai}'(z) \sim -z^{1/4} \exp(-\zeta) \sum_{r=0}^{m_3} c_r T_r^*(t)$$

$$\text{Bi}(z) \sim z^{-1/4} \exp(\zeta) \sum_{r=0}^{m_2} b_r T_r^*(t), \quad \text{Bi}'(z) \sim z^{1/4} \exp(\zeta) \sum_{r=0}^{m_4} d_r T_r^*(t)$$

where $t = (7/z)^{3/2}$.

In each case the degree of approximation, m_i , is such that the functions are determined accurate to at least nine significant figures if the function is greater than unity in magnitude and to nine decimal places if not. However, in case (a) the final accuracy is limited by the accuracy to which the sine and cosine of $\zeta + \pi/4$ can be computed. In the cases (a), (b), and (d) the Chebyshev coefficients were computed by rearrangement of the corresponding power series expansions using [3, p. 52, eq. 23]. In case (c) the coefficients were computed by numerical integration based on [3, p. 65, eq. 1], where the required Airy function values were computed in high precision using the Taylor expansions. In all four cases use is made of the Rice algorithm for summing the Chebyshev series (see [3, p. 57, eq. 50]).

Program

The subroutine is named AIRY and is composed of the four subroutines COEF1, COEF2, COEF3, and COEF4 corresponding, respectively, to the previous cases (a), (b), (c), and (d). The calling statement is

CALL AIRY (Z, NFUNC, ISCAL, AI, BI, AID, BID)

where the input variables Z, NFUNC, and ISCAL are defined as follows:

Z the prescribed argument

NFUNC an indicator for which of the functions are to be calculated, as follows:

$$\text{NFUNC} \begin{cases} \text{negative, for Ai' and Bi' only;} \\ 0, & \text{for all four functions;} \\ \text{positive, for Ai and Bi only} \end{cases}$$

ISCAL an indicator for whether or not in case (d) the following scaling is required:

$$(A_i, A_i') = (A_i, A_i') \exp(\zeta);$$

$$(B_i, B_i') = (B_i, B_i') \exp(-\zeta);$$

if ISCAL = 0 there is no scaling; otherwise, scaling as above.

The output variables will contain the respective values of A_i , B_i , A_i' , and B_i' . After the CALL statement the input parameters are left unchanged.

The relevant prescribed Chebyshev coefficients are stored in arrays A , B , C , and D corresponding, respectively, to a_r , b_r , c_r , and d_r in the previous analysis. Subroutine CHEB evaluates the sum of the required Chebyshev series given the argument x ($0 \leq x \leq 1$) and the corresponding N Chebyshev coefficients a_0, a_1, \dots, a_{N-1} stored as $A(I)$, $I = 1, N$.

Test Results

Tests have been carried out on an ICL 1905E computer (11S floating-point arithmetic). The Wronskian relationship $W = A_i B_i' - A_i' B_i = 1/\pi$ has been used to check on the results, as have the tables in [1] and the results obtained from a high precision calculation based on the Taylor expansions.

Tests have also been carried out for comparison with ACM Algorithm 301 [2], which gives at least eight-figure accuracy; the present algorithm is comparable in time for setting up large tables of the functions, and is certainly more efficient when only a small number of function evaluations are required outside the asymptotic ranges. For shorter word machines the user may truncate the given Chebyshev coefficients. The routine has run successfully on a Data General Corporation Nova 840 (7S floating-point arithmetic), with at least five-figure accuracy obtained.

REFERENCES

1. ABRAMOWITZ, M., AND STEGUN, I.A. *Handbook of Mathematical Functions*. Dover Publications, New York, 1965, pp. 446, 448-449, 475-477.
2. BOND, G., AND PITTEWAY, M.L.V. Algorithm 301, Airy function. *Comm. ACM* 10, 5 (May 1967), 291-293.
3. FOX, L., AND PARKER, I.B. *Chebyshev Polynomials in Numerical Analysis*. Oxford U. Press, London, 1968, pp. 48-58, 65-68.

ALGORITHM

SUBROUTINE AIRY(Z, NFUNC, ISCAL, AI, BI, AID, BID)	AIR 10
C P.J.PRINCE, MATHEMATICS DEPARTMENT, TEESIDE POLYTECHNIC,	AIR 20
C MIDDLESBROUGH, CLEVELAND TS1 3BA, GREAT BRITAIN.	AIR 30
C JANUARY 9, 1975 .	AIR 40
C REFERENCES -	AIR 50
C (1) - ABRAMOWITZ, M. AND STEGUN, I.A. HANDBOOK OF	AIR 60
C MATHEMATICAL FUNCTIONS. DOVER PUBLICATIONS INC.,	AIR 70
C NEW YORK, 1965 PAGES 446, 448-449, 475-477 .	AIR 80
C (2) - FOX, L. AND PARKER, I.B. CHEBYSHEV POLYNOMIALS IN	AIR 90
C NUMERICAL ANALYSIS. OXFORD UNIVERSITY PRESS, 1968	AIR 100
C PAGES 48-58, 65-68 .	AIR 110
C (3) - BOND, GILLIAN AND PITTEWAY, M.L.V. ALGORITHM 301,	AIR 120
C AIRY FUNCTION(S20). C.A.C.M. VOL. 10, NO. 5, MAY 1967	AIR 130
C THIS SUBROUTINE EVALUATES THE AIRY FUNCTIONS AND THEIR	AIR 140
C DERIVATIVES FOR ANY REAL ARGUMENT WITHIN COMPUTER	AIR 150
C CAPABILITY, BY MEANS OF CHEBYSHEV SERIES APPROXIMATIONS.	AIR 160
C Z - THE ARGUMENT FOR WHICH THE FUNCTIONS ARE	AIR 170
C TO BE EVALUATED.	AIR 180
C AI, BI - THE AIRY FUNCTIONS.	AIR 190
C AID, BID - THEIR RESPECTIVE DERIVATIVES.	AIR 200
C NFUNC - INDICATES WHICH OF THE FUNCTIONS ARE TO BE	AIR 210
C CALCULATED.	AIR 220
C ISCAL - INDICATES WHETHER OR NOT SCALING IS	AIR 230
C REQUIRED.	AIR 240
C FOUR RANGES OF ARGUMENT Z ARE CONSIDERED -	AIR 250
C (A) Z .LE. -7.0	AIR 260
C (B) -7.0 .LT. Z .LE. 0.0	AIR 270
C (C) 0.0 .LT. Z .LT. 7.0	AIR 280
C (D) Z .GE. 7.0	AIR 290
C THE SUBROUTINES COEF1, COEF2, COEF3 AND COEF4 WHICH ARE	AIR 300
C CALLED BY THIS SUBROUTINE CORRESPOND RESPECTIVELY TO	AIR 310
C THESE FOUR RANGES.	AIR 320
C THE FUNCTIONS CALCULATED ARE AS FOLLOWS -	AIR 330
C IF NFUNC IS NEGATIVE AID, BID.	AIR 340

```

C IF NFUNC = 0 ALL FOUR. AIR 350
C IF NFUNC IS POSITIVE AI,BI. AIR 360
C ISCAL INDICATES WHETHER OR NOT IN RANGE (D) THE FOLLOWING AIR 370
C SCALING IS TO BE CARRIED OUT - AIR 380
C (AI,AID) = (AI,AID)*EXP(ZETA), (BI,BID) = (BI,BID)*EXP(-ZETA) AIR 390
C WHERE ZETA = 2.0/3.0*Z**1.5 AIR 400
C IF ISCAL = 0 THEN THERE IS NO SCALING OTHERWISE SCALING AIR 410
C AS ABOVE. AIR 420
C IN (A) THE FINAL ACCURACY IS HEAVILY DEPENDENT ON THE AIR 430
C ACCURACY TO WHICH THE SINE AND COSINE OF AIR 440
C ANG = 2.0/3.0*(-Z)**1.5+PI/4 CAN BE COMPUTED. AIR 450
C IF ANGLM .LT. ANG .LE. ANGUP THEN SOME ACCURACY MAY BE AIR 460
C LOST. A NON FATAL ERROR WARNING IS GIVEN. AIR 470
C IF ANG .GT. ANGUP A NON FATAL ERROR WARNING IS GIVEN AND AIR 480
C THE FOUR FUNCTIONS ARE ASSIGNED THE VALUE ZERO. AIR 490
C IN (D) IF ISCAL = 0 OVERFLOW IN EXP WILL OCCUR IF AIR 500
C Z .GE. ZLIM. IN THIS CASE A NON FATAL ERROR WARNING IS AIR 510
C GIVEN AND THE FOLLOWING VALUES ARE ASSIGNED - AIR 520
C AI = AID = 0.0 AIR 530
C BI = BID = EXP(ZLIM). AIR 540
C IF ISCAL .NE. 0 OVERFLOW WILL OCCUR IF AIR 550
C Z .GT. ZUP = ZMAX**(2.0/3.0) WHERE ZMAX IS THE MAXIMUM AIR 560
C REPRESENTABLE NUMBER ON THE COMPUTER. IN THIS CASE A AIR 570
C NON FATAL ERROR WARNING IS GIVEN AND THE FOUR FUNCTIONS AIR 580
C ARE ASSIGNED THE VALUE ZERO. AIR 590
C NOUT IS THE OUTPUT CHANNEL USED. AIR 600
C THE MACHINE DEPENDENT CONSTANTS (VALUES BEING ASSIGNED BY AIR 610
C APPROPRIATE DATA STATEMENTS) ARE - AIR 620
C IN COEF1 - ANGLM,ANGUP AND NOUT. AIR 630
C IN COEF4 - ZLIM,ZUP AND NOUT. AIR 640
C FOR I.C.L. 1905E(11S FLOATING POINT ARITHMETIC) - AIR 650
C ANGLM = 250.0, ANGUP = 1.0E10, NOUT = 2, ZLIM = 175.0, AIR 660
C ZUP = 1.0E50, ZMAX = 5.6E76 AIR 670
C IF (Z) 10, 10, 40 AIR 680
C 10 IF (Z+7.0) 20, 20, 30 AIR 690
C 20 CALL COEF1(Z, NFUNC, AI, BI, AID, BID) AIR 700
C RETURN AIR 710
C 30 CALL COEF2(Z, NFUNC, AI, BI, AID, BID) AIR 720
C RETURN AIR 730
C 40 IF (Z-7.0) 50, 60, 60 AIR 740
C 50 CALL COEF3(Z, NFUNC, AI, BI, AID, BID) AIR 750
C RETURN AIR 760
C 60 CALL COEF4(Z, NFUNC, ISCAL, AI, BI, AID, BID) AIR 770
C RETURN AIR 780
C END AIR 790

SUBROUTINE COEF1(Z, NFUNC, AI, BI, AID, BID) COE 10
C THIS SUBROUTINE EVALUATES THE FOUR FUNCTIONS IN THE CASE COE 20
C WHEN Z .LE. -7.0 USING CHEBYSHEV SERIES APPROXIMATIONS TO COE 30
C THE CORRESPONDING ASYMPTOTIC EXPANSIONS. COE 40
C DIMENSION A(5), B(5), C(5), D(5) COE 50
C DATA A(1), A(2), A(3), A(4), A(5) /1.1282427601, COE 60
C * -0.6803534E-04,0.16687E-06,-0.128E-08,0.2E-10/ COE 70
C DATA B(1), B(2), B(3), B(4), B(5) /0.7822108673E-01, COE 80
C * -0.6895649E-04,0.32857E-06,-0.370E-08,0.7E-10/ COE 90
C DATA C(1), C(2), C(3), C(4), C(5) /1.1285404716,0.8046925E-04, COE 100
C * -0.18161E-06,0.135E-08,-0.2E-10/ COE 110
C DATA D(1), D(2), D(3), D(4), D(5) /-0.10954855184, COE 120
C * 0.7713350E-04,-0.35168E-06,0.388E-08,-0.7E-10/ COE 130
C DATA PI4 /0.78539816340/ COE 140
C DATA ANGLM, ANGUP /250.0,1.0E10/ COE 150
C DATA NOUT /2/ COE 160
C X = -Z COE 170
C SX = SQRT(X) COE 180
C Y = X*SX COE 190
C ZETA = 0.66666666667*Y COE 200
C Z4 = SQRT(SX) COE 210
C ANG = ZETA + PI4 COE 220
C TEST ARGUMENT SIZE FOR SIN AND COS. COE 230
C IF (ANGLM-ANG) 10, 40, 40 COE 240
C 10 IF (ANGUP-ANG) 20, 20, 30 COE 250
C 20 WRITE (NOUT,99999) Z, ANG COE 260
C AI = 0.0 COE 270
C BI = 0.0 COE 280
C AID = 0.0 COE 290
C BID = 0.0 COE 300
C RETURN COE 310
C 30 WRITE (NOUT,99998) Z, ANG COE 320
C 40 SN = SIN(ANG) COE 330
C CN = COS(ANG) COE 340
C ZETA1 = 1.0/ZETA COE 350
C X = 7.0/X COE 360
C X = X*X*X COE 370
C ARGUMENT SCALED TO RANGE (0,1). COE 380
C EVALUATE THE RELEVANT SERIES. COE 390
C IF (NFUNC.LT.0) GO TO 50 COE 400
C CALL CHEB(X, 5, A, FA) COE 410
C CALL CHEB(X, 5, B, FB) COE 420
C Z4I = 1.0/Z4 COE 430

```

```

      FB = FB*ZETA1
      AI = Z4I*(SN*FA-CN*FB)
      BI = Z4I*(CN*FA+SN*FB)
      IF (NFUNC.GT.0) RETURN
50  CALL CHEB(X, 5, C, FC)
      CALL CHEB(X, 5, D, FD)
      FD = FD*ZETA1
      AID = -Z4*(CN*FC+SN*FD)
      BID = Z4*(SN*FC-CN*FD)
      RETURN
99999 FORMAT (//38X, 3H***//5X, 28HANGLE OUTSIDE MACHINE RANGE,,
* 38H THE FOUR FUNCTIONS HAVE BEEN ASSIGNED/5X, 11HTHE VALUE Z,
* 4HERO.//5X, 4HZ = , E20.8, 10H ANGLE = , E20.8//38X, 3H***//
* )
99998 FORMAT (//38X, 3H***//5X, 25HSOME ACCURACY MAY BE LOST,
* 16H IN SIN AND COS.//5X, 4HZ = , E20.8, 10H ANGLE = ,
* E20.8//38X, 3H***//)
      END
COE 440
COE 450
COE 460
COE 470
COE 480
COE 490
COE 500
COE 510
COE 520
COE 530
COE 540
COE 550
COE 560
COE 570
COE 580
COE 590
COE 600
COE 610

      SUBROUTINE COEF2(Z, NFUNC, AI, BI, AID, BID)
      C THIS SUBROUTINE EVALUATES THE FOUR FUNCTIONS IN THE CASE
      C WHEN -7.0 .LT. Z .LE. 0.0 USING CHEBYSHEV SERIES
      C APPROXIMATIONS TO THE CORRESPONDING TAYLOR EXPANSIONS.
      DIMENSION A(17), B(16), C(16), D(17)
      DATA A(1), A(2), A(3), A(4), A(5), A(6), A(7), A(8), A(9),
* A(10), A(11), A(12), A(13), A(14), A(15), A(16), A(17)
* /0.11535880704,0.6542816649E-01,0.26091774326,0.21959346500,
* 0.12476382168,-0.43476292594,0.28357718605,-0.9751797082E-01,
* 0.2182551823E-01,-0.350454097E-02,0.42778312E-03,
* -0.4127564E-04,0.323880E-05,-0.21123E-06,0.1165E-07,
* -0.55E-09,0.2E-10/
      DATA B(1), B(2), B(3), B(4), B(5), B(6), B(7), B(8), B(9),
* B(10), B(11), B(12), B(13), B(14), B(15), B(16)
* /0.10888288487,-0.17511655051,0.13887871101,-0.11469998998,
* 0.22377807641,-0.18546243714,0.8063565186E-01,
* -0.2208647864E-01,0.422444527E-02,-0.60131028E-03,
* 0.6653890E-04,-0.590842E-05,0.43131E-06,-0.2638E-07,
* 0.137E-08,-0.6E-10/
      DATA C(1), C(2), C(3), C(4), C(5), C(6), C(7), C(8), C(9),
* C(10), C(11), C(12), C(13), C(14), C(15), C(16)
* /0.7571648463E-01,-0.10150232871,0.7800551669E-01,
* -0.8324569361E-01,0.10105322731,-0.6578603344E-01,
* 0.2500140353E-01,-0.625962704E-02,0.111945149E-02,
* -0.15102718E-03,0.1598086E-04,-0.136545E-05,0.9636E-07,
* -0.572E-08,0.29E-09,-0.1E-10/
      DATA D(1), D(2), D(3), D(4), D(5), D(6), D(7), D(8), D(9),
* D(10), D(11), D(12), D(13), D(14), D(15), D(16), D(17)
* /0.61603048107,0.85738069722,0.86345334421,0.80890228699,
* -0.50565665369,-0.81829752697,0.77829538563,-0.31201242692,
* 0.7677186198E-01,-0.1320520264E-01,0.170185698E-02,
* -0.17177956E-03,0.1401068E-04,-0.94532E-06,0.5374E-07,
* -0.261E-08,0.11E-09/
      DATA E1, E2, ROOT3 /0.35502805389,0.25881940379,1.7320508076/
      X = -0.29154518950E-02*Z*Z*Z
      C ARGUMENT SCALED TO RANGE (0,1).
      C EVALUATE THE RELEVANT SERIES.
      IF (NFUNC.LT.0) GO TO 10
      CALL CHEB(X, 17, A, FA)
      CALL CHEB(X, 16, B, FB)
      FA = E1*FA
      FB = E2*Z*FB
      AI = FA - FB
      BI = ROOT3*(FA+FB)
      IF (NFUNC.GT.0) RETURN
10  CALL CHEB(X, 16, C, FC)
      CALL CHEB(X, 17, D, FD)
      FC = E1*Z*Z*FC
      FD = E2*FD
      AID = FC - FD
      BID = ROOT3*(FC+FD)
      RETURN
      END
COE 10
COE 20
COE 30
COE 40
COE 50
COE 60
COE 70
COE 80
COE 90
COE 100
COE 110
COE 120
COE 130
COE 140
COE 150
COE 160
COE 170
COE 180
COE 190
COE 200
COE 210
COE 220
COE 230
COE 240
COE 250
COE 260
COE 270
COE 280
COE 290
COE 300
COE 310
COE 320
COE 330
COE 340
COE 350
COE 360
COE 370
COE 380
COE 390
COE 400
COE 410
COE 420
COE 430
COE 440
COE 450
COE 460
COE 470
COE 480
COE 490
COE 500
COE 510
COE 520
COE 530

      SUBROUTINE COEF3(Z, NFUNC, AI, BI, AID, BID)
      C THIS SUBROUTINE EVALUATES THE FOUR FUNCTIONS IN THE CASE
      C WHEN 0.0 .LT. Z .LT. 7.0 USING CHEBYSHEV SERIES
      C APPROXIMATIONS TO THE CORRESPONDING WEIGHTED FUNCTIONS.
      DIMENSION A(20), B(25), C(22), D(24)
      DATA A(1), A(2), A(3), A(4), A(5), A(6), A(7), A(8), A(9),
* A(10), A(11), A(12), A(13), A(14), A(15), A(16), A(17),
* A(18), A(19), A(20) /1.2974695794,-0.20230907821,
* -0.45786169516,0.12953331987,0.6983827954E-01,
* -0.3005148746E-01,-0.493036334E-02,0.390425474E-02,
* -0.1546539E-04,-0.32193999E-03,0.3812734E-04,0.1714935E-04,
* -0.416096E-05,-0.50623E-06,0.26346E-06,-0.281E-08,
* -0.1122E-07,0.120E-08,0.31E-09,-0.7E-10/
      DATA B(1), B(2), B(3), B(4), B(5), B(6), B(7), B(8), B(9),
COE 10
COE 20
COE 30
COE 40
COE 50
COE 60
COE 70
COE 80
COE 90
COE 100
COE 110
COE 120
COE 130
COE 140

```

```

* B(10), B(11), B(12), B(13), B(14), B(15), B(16), B(17), COE 150
* B(18), B(19), B(20), B(21), B(22), B(23), B(24), B(25) COE 160
* /0.47839902387,-0.6881732880E-01,0.20936146768, COE 170
* -0.3988095886E-01,0.4758441683E-01,-0.812296149E-02, COE 180
* 0.462845913E-02,0.70010098E-03,-0.75611274E-03, COE 190
* 0.68958657E-03,-0.33621865E-03,0.14501663E-03,-0.4766359E-04, COE 200
* 0.1251965E-04,-0.193012E-05,-0.19032E-06,0.29390E-06, COE 210
* -0.13436E-06,0.4231E-07,-0.967E-08,0.135E-08,0.7E-10, COE 220
* -0.12E-09,0.4E-10,-0.1E-10/ COE 230
DATA C(1), C(2), C(3), C(4), C(5), C(6), C(7), C(8), C(9), COE 240
* C(10), C(11), C(12), C(13), C(14), C(15), C(16), C(17), COE 250
* C(18), C(19), C(20), C(21), C(22) /-2.2359158747, COE 260
* -0.2638272392E-01,0.95151904332,-0.8383641182E-01, COE 270
* -0.19401303219,0.3580664778E-01,0.2269348562E-01, COE 280
* -0.671179820E-02,-0.152460473E-02,0.75474150E-03, COE 290
* 0.3729934E-04,-0.5653536E-04,0.350796E-05,0.289418E-05, COE 300
* -0.47423E-06,-0.9449E-07,0.3054E-07,0.109E-08,-0.130E-08, COE 310
* 0.8E-10,0.4E-10,-0.1E-10/ COE 320
DATA D(1), D(2), D(3), D(4), D(5), D(6), D(7), D(8), D(9), COE 330
* D(10), D(11), D(12), D(13), D(14), D(15), D(16), D(17), COE 340
* D(18), D(19), D(20), D(21), D(22), D(23), D(24) COE 350
* /0.71364662990,0.23777925892,0.28219009446,0.4912480040E-01, COE 360
* 0.6741261353E-01,-0.406308553E-02,0.1544814895E-01, COE 370
* -0.449172894E-02,0.322474426E-02,-0.105361380E-02, COE 380
* 0.41311371E-03,-0.8536169E-04,0.655166E-05,0.960458E-05, COE 390
* -0.641792E-05,0.280308E-05,-0.89454E-06,0.21392E-06, COE 400
* -0.2958E-07,-0.309E-08,0.376E-08,-0.148E-08,0.39E-09, COE 410
* -0.7E-10/ COE 420
X = 0.14285714286*X COE 430
EX = EXP(1.75*X) COE 440
EY = 1.0/EX COE 450
C ARGUMENT SCALED TO RANGE (0,1). COE 460
C EVALUATE THE RELEVANT SERIES. COE 470
IF (NFUNC.LT.0) GO TO 10 COE 480
CALL CHEB(X, 20, A, AI) COE 490
CALL CHEB(X, 25, B, BI) COE 500
C MULTIPLY BY APPROPRIATE WEIGHTING FACTORS. COE 510
AI = AI*EY COE 520
BI = BI*EX COE 530
IF (NFUNC.GT.0) RETURN COE 540
10 CALL CHEB(X, 22, C, AID) COE 550
CALL CHEB(X, 24, D, BID) COE 560
C MULTIPLY BY APPROPRIATE WEIGHTING FACTORS. COE 570
AID = AID*EY COE 580
BID = BID*EX COE 590
RETURN COE 600
END COE 610

SUBROUTINE COEF4(Z, NFUNC, ISCAL, AI, BI, AID, BID) COE 10
C THIS SUBROUTINE EVALUATES THE FOUR FUNCTIONS (SCALED COE 20
C UNLESS ISCAL = 0) IN THE CASE WHEN Z .GE. 7.0 USING COE 30
C CHEBYSHEV SERIES APPROXIMATIONS TO THE CORRESPONDING COE 40
C ASYMPTOTIC EXPANSIONS. COE 50
DIMENSION A(7), B(7), C(7), D(7) COE 60
DATA A(1), A(2), A(3), A(4), A(5), A(6), A(7) COE 70
* /0.56265126169,-0.76136219E-03,0.765252E-05,-0.14228E-06, COE 80
* 0.380E-08,-0.13E-09,0.1E-10/ COE 90
DATA B(1), B(2), B(3), B(4), B(5), B(6), B(7) COE 100
* /1.1316635302,0.166141673E-02,0.1968882E-04,0.47047E-06. COE 110
* 0.1769E-07,0.94E-09,0.6E-10/ COE 120
DATA C(1), C(2), C(3), C(4), C(5), C(6), C(7) COE 130
* /0.56635357764,0.107273242E-02,-0.910034E-05,0.15998E-06, COE 140
* -0.415E-08,0.14E-09,-0.1E-10/ COE 150
DATA D(1), D(2), D(3), D(4), D(5), D(6), D(7) COE 160
* /1.1238058432,-0.230925296E-02,-0.2309457E-04,-0.52171E-06. COE 170
* -0.1907E-07,-0.100E-08,-0.7E-10/ COE 180
DATA ZLIM, ZUP /175.0,1.0E50/ COE 190
DATA NOUT /2/ COE 200
C TEST FOR OVERFLOW IN Z**1.5 . COE 210
IF (ZUP-Z) 10, 20, 20 COE 220
10 WRITE (NOUT,99999) Z COE 230
AI = 0.0 COE 240
BI = 0.0 COE 250
AID = 0.0 COE 260
BID = 0.0 COE 270
RETURN COE 280
20 SZ = SQRT(Z) COE 290
Y = Z*SZ COE 300
Z4 = SQRT(SZ) COE 310
IF (ISCAL) 30, 40, 30 COE 320
30 EX = 1.0 COE 330
EY = 1.0 COE 340
GO TO 70 COE 350
40 ZETA = 0.66666666667*Y COE 360
C TEST FOR OVERFLOW IN EXP COE 370
IF (ZLIM-ZETA) 50, 60, 60 COE 380
50 WRITE (NOUT,99998) Z, ZETA COE 390
AI = 0.0 COE 400
AID = 0.0 COE 410

```



```

      BI = EXP(ZLIM)
      BID = BI
      RETURN
60 EX = EXP(ZETA)
   EY = 1.0/EX
70 X = 18.520259177/Y
C ARGUMENT SCALED TO RANGE (0,1).
C EVALUATE THE RELEVANT SERIES.
   IF (NFUNC.LT.0) GO TO 80
   CALL CHEB(X, 7, A, AI)
   CALL CHEB(X, 7, B, BI)
   Z4I = 1.0/24
   AI = Z4I*AI*EY
   BI = Z4I*BI*EX
   IF (NFUNC.GT.0) RETURN
80 CALL CHEB(X, 7, C, AID)
   CALL CHEB(X, 7, D, BID)
   AID = -Z4*AID*EY
   BID = Z4*BID*EX
   RETURN
99999 FORMAT (//38X, 3H**//5X, 18HOVERFLOW IN Z**1.5//5X, 4HZ = ,
* E20.8//5X, 41HTHE FOUR FUNCTIONS HAVE BEEN ASSIGNED THE,
* 12H VALUE ZERO.//38X, 3H**//)
99998 FORMAT (//38X, 3H**//5X, 21HOVERFLOW IN EXP(ZETA)//5X,
* 4HZ = , E20.8, 9H ZETA = , E20.8//5X, 18HTHE FOLLOWING VALU,
* 23HES HAVE BEEN ASSIGNED -//5X, 25HAI = AID = 0.0, BI = BID ,
* 11H= EXP(ZLIM)//38X, 3H**//)
END

```

COE 420
COE 430
COE 440
COE 450
COE 460
COE 470
COE 480
COE 490
COE 500
COE 510
COE 520
COE 530
COE 540
COE 550
COE 560
COE 570
COE 580
COE 590
COE 600
COE 610
COE 620
COE 630
COE 640
COE 650
COE 660
COE 670
COE 680
COE 690

```

      SUBROUTINE CHEB(X, N, A, F)
C THIS SUBROUTINE APPLIES THE RICE ALGORITHM TO THE SUM
C F(X) = SIGMA-PRIME (R=1 TO N) A(R)*TSTAR-SUB-(R-1)(X)
C WHERE THE TSTAR-SUB-R(X) ARE THE SHIFTED CHEBYSHEV
C POLYNOMIALS.
C X - THE ARGUMENT FOR WHICH THE SERIES IS TO BE
C EVALUATED. (0.0 .LE. X .LE. 1.0).
C N - THE NUMBER OF TERMS OF THE SERIES TO BE SUMMED.
C A(R) - THE CHEBYSHEV COEFFICIENTS FOR THE SERIES.
C F - THE COMPUTED SUM OF THE SERIES.
      DIMENSION A(25)
      B = 0.0
      D = A(N)
      U = X + X - 1.0
      Y = U + U
      J = N - 1
      DO 10 I=3,N
         C = B
         B = D
         D = Y*B - C + A(J)
         J = J - 1
10 CONTINUE
      F = U*D - B + 0.5*A(1)
      RETURN
      END

```

CHE 10
CHE 20
CHE 30
CHE 40
CHE 50
CHE 60
CHE 70
CHE 80
CHE 90
CHE 100
CHE 110
CHE 120
CHE 130
CHE 140
CHE 150
CHE 160
CHE 170
CHE 180
CHE 190
CHE 200
CHE 210
CHE 220
CHE 230
CHE 240
CHE 250

REMARK ON ALGORITHM 498

Airy Functions Using Chebyshev Series Approximations [P.J. Price, *ACM Trans. Math. Softw.* 1, 4 (Dec. 1975), 372-379]

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Algorithm 498 presents a method for approximating the Airy functions and their first derivatives based on Chebyshev expansions. The method employed and the actual Chebyshev expansions presented give adequate routines for accuracies limited to about 8D (absolute in the negative region and relative in the positive). Such limited accuracy is adequate for single precision on 32- or 36-bit machines, that is, working precisions of 6D to 8D, approximately. Machines currently in common use regularly employ floating-point precisions of 11, 15, 17, and 18D. To cover this accuracy range and to deal with possible double-precision implementations, we have extended the accuracy capabilities of this basic algorithm to 30D

and new Chebyshev expansions have been generated to that precision [1]. This necessitated some modifications to the basic method of Algorithm 498. The most significant of such modifications was the need to introduce an additional range subdivision for the positive region.

Following as closely as possible the notation of Algorithm 498, we use for regions (a) and (b) entirely similar expansions, with the essential differences that the break point is chosen at -5 rather than -7 and we absorb the constants c_1 and c_2 of region (b) into the Chebyshev series. The former change is merely to improve the balance of the approximations for the important accuracy range 15D to 20D and to speed up the routines for the important argument range that includes the first few oscillations of the functions.

In order to increase the precision in the positive region it becomes necessary to introduce an additional region (c'), analogous to region (c). For the latter region we employ

$$\begin{aligned} Ai(z)\exp(+\alpha z) &= \sum_{r=0}' a_r T_r^*(t), & Ai'(z)\exp(+\beta z) &= \sum_{r=0}' c_r T_r^*(t) \\ Bi(z)\exp(-\beta z) &= \sum_{r=0}' b_r T_r^*(t), & Bi'(z)\exp(-\alpha z) &= \sum_{r=0}' d_r T_r^*(t) \end{aligned}$$

with $0 < z < 4.5$, $t = z/4.5$, $\alpha = 1.5$, and $\beta = 1.375$. The prime means that the term with suffix $r = 0$ in each expansion is halved. The constants α and β were chosen to approximately minimize the variation of the function being expanded as a polynomial, this being required to give good relative error control for the resulting approximations [2]. These constants were also chosen to be exactly representable on any existing machine, thus aiding portability and preserving accuracy.

For the additional region (c') we use

$$\begin{aligned} Ai(z)\exp(+\alpha z) &= \sum_{r=0}' a_r T_r^*(t), & Ai'(z)\exp(+\alpha z) &= \sum_{r=0}' c_r T_r^*(t) \\ Bi(z)\exp(-\beta z) &= \sum_{r=0}' b_r T_r^*(t), & Bi'(z)\exp(-\beta z) &= \sum_{r=0}' d_r T_r^*(t) \end{aligned}$$

with $4.5 < z < 9$, $t = (z - 4.5)/4.5$, $\alpha = 2.5$, and $\beta = 2.625$.

For region (d) we use expansions as in Algorithm 498, except that the break point of 9 rather than 7 is used. This allows both slightly faster convergence for the expansions and a more efficient and accurate calculation of t to be performed. The parameter t can now be calculated as

$$t = (9/z)^{3/2} = 27/(z)^{3/2} = 18/\zeta.$$

Thus we avoid the necessity of calculating a $3/2$ power more than once and/or of carrying a further nonexact constant $(7)^{3/2}$.

The modifications to the code of Algorithm 498 to employ these new expansions for obtaining a higher precision routine are fairly straightforward and the actual expansions have been published elsewhere [1]. Using these expansions the authors have produced routines that have been included in the S17 chapter of the NAG library [3]. The routines have been tested and have been shown to provide approximations accurate to within the errors expected for functions of this type for machines of precisions up to 18D, the errors being of the order of the machine precision except where losses are unavoidable. For instance, for large arguments where the errors are dominated by the imprecisions of sine, cosine, and exponential routines.

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ALGORITHM 499

An Efficient Scanning Technique [Z]

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Key Words and Phrases: array scanning, pattern recognition, partial differential equations, finite-differences, Laplace's equation

CR Categories: 4.22, 5.14, 5.17, 8.2

Language: Fortran

DESCRIPTION

1. Introduction

Scanning an array refers to selecting all the elements of the array in a particular order. We shall consider two scanning techniques: a raster scan and a new contour scan. The raster scan selects the elements of a row from left to right and rows sequentially from top to bottom. The contour scan selects the elements of an array so that they are adjacent to or lie on a given boundary and so that the subsequent elements are adjacent to the previously selected points. Use of the contour scanning technique, rather than the raster scanning technique, accelerates the speed of propagation of the boundary condition effects when solving partial differential equations by finite-difference methods or other iterative methods. This technique is useful in other applications, such as pattern recognition, although the implementation discussed here refers to iterative solutions of elliptic boundary value problems.

2. Basic Concept

Let R be an open rectangular two-dimensional region with a boundary C containing an open region R_1 with boundaries C_1 and C_2 as shown in Figure 1, and let the region R be discretized with S mesh points such that $S \subset R$. The algorithm described here selects all the points S_1 in a particular order such that $S_1 \subset R_0$ where $R_0 = R_1 \cup C_2$.

The selection of the points in R_0 starts from a point adjacent to a boundary point on C_1 . Then all the points in R_0 adjacent to C_1 and on C_2 are selected in a sequence as they occur in the clockwise direction of a path along C_1 and on C_2 . It is noted that line segments joining all the selected adjacent points form a polygon even if the boundary C_1 and C_2 do not constitute a closed contour and intersect the boundary C . Subsequently selected points in R_0 are adjacent to the previously selected points. The procedure is repeated until the last element of S_1 is found.

Special treatment is given to multiply connected boundaries.

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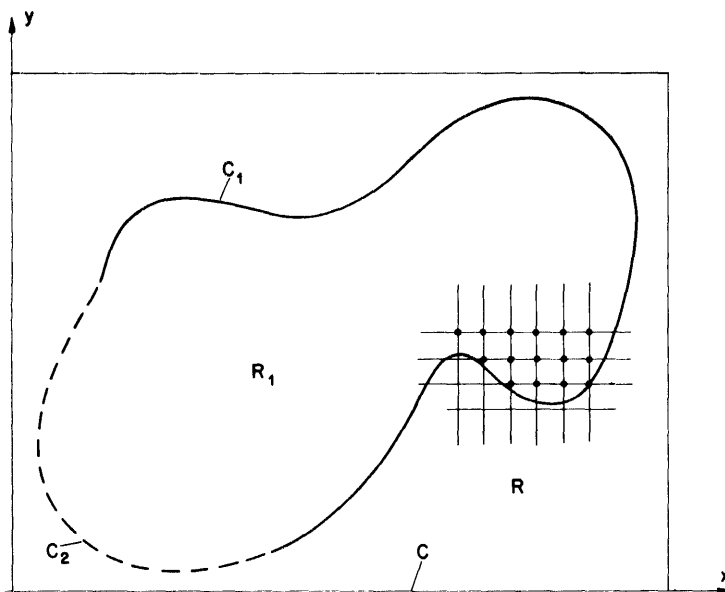


Fig. 1. Mixed nonconvex boundary value problem

3. Implementation

The simplest scanning procedure is the raster scan since it can be implemented with two DO loops. However, this procedure has two drawbacks: less favored boundary conditions affect the solution at the very end of the scan; and for nonrectangular boundaries the raster scan samples points outside the specified region where the solution is not desired.

A scanning technique with neither drawback is the contour scan shown in Figure 2. After the scan has been initialized by the proper selection of a starting point and an initial direction, it proceeds by trying to make a left turn in a two-dimensional array. If it encounters either a point on a Dirichlet boundary or a previously scanned point, it then tries to go straight ahead. If it fails, a raster scan is used to find the next available point within the boundaries. The raster scan terminates the procedure when all the points in the array have been selected.

It is seen from Figure 2 that the complexity of the path traced by the contour scan depends on the choice of the starting point. This starting point may be specified by the user. However, a more efficient way of choosing a proper starting point and the optimum contour path is guaranteed by the subroutine CONOPT. The subroutine uses two criteria for selecting proper starting points: first, it tries to choose a point adjacent to a mesh point with the highest value of Dirichlet boundary condition; second, it tries to select a point so that the contour generated is a continuous curve. This latter criterion tends to produce a consistent ordering [2].

The subroutine HIVA locates the points with large numerical values of Dirichlet boundaries by means of the raster scan. The subroutine SPOINT then selects up to twenty choices for possible starting points. The points are chosen in such a way as to minimize the number of corners in the scan and so that the scan runs adjacent, if possible, to the entire equipotential boundary. The subroutine CONTR generates different orderings from each of these points. The subroutine CONOPT chooses either the first continuous scan, or the scan with the fewest discontinuities.

The initial direction in which the contour scan proceeds from a starting point is determined automatically in the subroutine CONTR as follows. If the Dirichlet boundary point is located at $(I - 1, J)$ and the point, $(I, J + 1)$ is an interior "free" point, the initial direction indicator, K , is set to 1, and the turn indicator, L , is set to 2, i.e. the scanning proceeds in the direction of increasing J . If this condition is not met, the previous pattern is rotated about the point (I, J) . Eventually, one of four possible directions will be chosen.

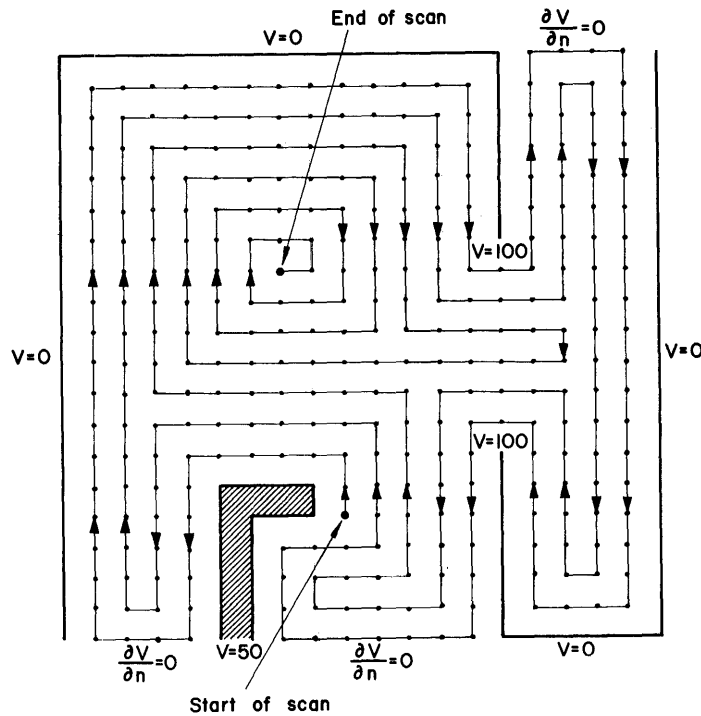


Fig. 2. Sample of boundary conditions and contour scan

The scanning path is computed only once for each different set of Dirichlet boundary values and it is stored in a tabular form. The scanning then is carried out by a single DO loop using table look-up. One can change the direction of the scan by altering the table look-up. For repeated solutions of the same problem one may punch out the selected contour scan.

4. Input-Output Considerations

In order to determine a contour scanning path, data is required in the NET array locating Dirichlet boundaries. The optimization of the contour scan requires numerical values of these boundaries. If the boundaries are piecewise linear or are described by a simple function, they may be generated by the computer. Otherwise the user has to supply them point by point. The numerical values of the boundaries are introduced in the array A and the corresponding location flags are automatically shifted in the array NET since the dimension of the array NET is greater than that of the array A. The program can handle purely Neumann problems, i.e. with the boundary consisting of C_2 only. Setting the boundary conditions in the arrays A and NET may be implemented in many different ways and it is felt that it does not belong to the described algorithm.

The intermediate results printed by the subprogram CONOPT may be suppressed by setting the flag IPRINT to zero.

5. Efficiency of the Algorithm

This algorithm has been designed to be used in iterative solutions of elliptic boundary-value problems. Therefore, its efficiency can be measured by comparing the total machine time spent for solving a given problem when using the contour scan and the raster scan. We consider two stages: selecting and storing a set of points, S_1 , from the region $R_1 \cup C_2$; and solving the problem by iteration.

Since the implementation of the contour scan is more complex than the implementation of the raster scan, the first stage is more economical for the raster scan, depending on the geometry of the region R_0 . Various test problems indicate that the set-up time for the contour scan is 10 to 50 times longer than that for the raster scan. However, this set-up time is only 10 to .5 times longer than the average

machine time per iteration for a small problem, as in Figure 2, and for large problems with up to 6500 points, respectively.

The average machine time per iteration during the second stage is exactly the same for both scanning techniques since the number of elements in S_1 is the same in both cases. It has been found experimentally [1] that the ordering of S_1 generating the contour scan causes a higher rate of convergence of the finite-difference methods as compared to the ordering generated by the raster scan. In all the test problems, this faster convergence resulted in 2 to 10 percent fewer iterations when using the Jacobi, Gauss-Seidel, or SOR methods. A more efficient method [1] resulted in up to 60 percent less iterations. Other scanning techniques such as the diagonal scan and the "red-black" scan [2] give the same rate of convergence as the raster scan.

The subroutine CONTR may be used as an independent subprogram for generating only one contour scanning path. In this case, it requires only 2 to 5 times more set-up time than the raster scan.

All the test problems have been run on a CDC-6400 computer.

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ALGORITHM

```

SUBROUTINE CONOPT(M, N, MM, NN, MN, A, NET, IPRINT, NP, IJ)      CON  10
C THIS SUBPROGRAM OPTIMIZES THE CONTOUR SCAN AND IT RETURNS    CON  20
C THE OPTIMUM CONTOUR SCANNING PATH STORED IN THE ARRAY IJ.   CON  30
C      ....INPUT VARIABLES....                                CON  40
C M -NUMBER OF SUBDIVISIONS IN X-DIRECTION. DIMENSION OF THE CON  50
C   ARRAY A IN I DIRECTION.                                  CON  60
C N -AS ABOVE BUT IN Y AND J DIRECTIONS.                    CON  70
C A -WORKING ARRAY WITH NUMERICAL VALUES OF DIRICHLET       CON  80
C   BOUNDARY CONDITIONS. A(M,N).                            CON  90
C MM -NUMBER OF SUBDIVISIONS IN X-DIRECTION. DIMENSION OF THE CON 100
C   ARRAY NET IN I DIRECTION. MM=M+2.                        CON 110
C NN -AS ABOVE BUT IN Y AND J DIRECTIONS. NN=N+2.          CON 120
C NET-AN INTEGER ARRAY WITH FLAGS LOCATING THE DIRICHLET     CON 130
C   AND NEUMANN BOUNDARY CONDITIONS. 0=VARIABLES,           CON 140
C   1=DIRICHLET B.C., 2=NEUMANN B.C. NET(MM,NN). THESE     CON 150
C   FLAGS ARE SHIFTED WITH RESPECT TO THE ARRAY A SUCH     CON 160
C   THAT THE LOCATION OF A POINT ON THE DIRICHLET BOUNDARY CON 170
C   A(I,J) CORRESPONDS TO NET(I+1,J+1).                     CON 180
C MN -DIMENSION OF THE ARRAY IJ. IT DEPENDS ON THE LENGTH   CON 190
C   OF THE CONTOUR PATH, NR. PREFERABLY MN=M*N.             CON 200
C IPRINT -IF IPRINT=0, THE PRINTOUT IN CONOPT IS SUPPRESSED. CON 210
C      ....OUTPUT VARIABLES....                               CON 220
C NP -NUMBER OF POINTS SCANNED BY THE CONTOUR PATH DETERMINED CON 230
C   BY THE SUBROUTINE CONTR.                                 CON 240
C IJ -AN INTEGER ARRAY CONTAINING COORDINATES I,J OF THE     CON 250
C   POINTS (I,J) SCANNED BY THE CONTOUR PATH. THEY ARE     CON 260
C   DETERMINED BY CONTR. THE ARRAY IJ HAS 3 COLUMNS, THE   CON 270
C   LAST ONE IS USED FOR COMPUTATIONAL PURPOSES IN CONTR,   CON 280
C   AND, ON THE OUTPUT, IT CONTAINS THE CORRESPONDING FLAGS CON 290
C   FROM THE ARRAY NET.                                     CON 300
C      ....INTERNAL VARIABLES....                             CON 310
C K -NUMBER OF STARTING POINTS DETERMINED BY THE SUBROUTINE  CON 320
C   SPOINT. K IS USED ONLY IN THIS SUBROUTINE.              CON 330
C IJS-AN INTEGER ARRAY. IT CONTAINS 4 COLUMNS, THE FIRST TWO CON 340
C   OF THEM CONTAIN COORDINATES I,J OF THE STARTING POINTS CON 350
C   (I,J) DETERMINED BY THE SUBROUTINE SPOINT, THE THIRD    CON 360
C   ONE CONTAINS CODED VALUES OF DIRICHLET B.C. ADJACENT   CON 370
C   TO A POINT (I,J), AND THE LAST ONE CONTAINS THE NUMBER  CON 380
C   OF BREAKS IN THE CONTOUR SCAN WHICH IS INITIATED FROM  CON 390
C   THE CORRESPONDING POINT (I,J). THE NUMBER OF ROWS IN   CON 400
C   IJS CAN BE CHANGED, IF NECESSARY.                       CON 410
C KK -NUMBER OF HIGH VALUE DIRICHLET BOUNDARY CONDITIONS,   CON 420
C   DETERMINED BY THE SUBROUTINE HIVA. KK IS USED IN SPOINT CON 430
C AA -AN ARRAY CONTAINING THE HIGH VALUE DIRICHLET B.C.,    CON 440
C   FOUND BY HIVA. THE LENGTH OF THE ARRAY AA CAN BE       CON 450
C   CHANGED, IF NECESSARY.                                  CON 460
C      ....NOTES....                                         CON 470
C THE PARAMETERS M,N,MM,NN,A,NET AND IPRINT MUST BE SUPPLIED CON 480
C BY THE USER (INPUT). PARAMETERS NP AND IJ ARE RETURNED BY CON 490
C THE SUBROUTINE CONOPT (OUTPUT). THE MAIN PROGRAM COMMUNI- CON 500
C CATES WITH THE SUBPROGRAM CONOPT THROUGH THE PARAMETER LIST CON 510
C ONLY. ALL THE LABELED COMMON STORAGE WITHIN THE SUBPROG-  CON 520
C RAM CONOPT IS FOR THE INTERNAL USE ONLY.                  CON 530

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C IF THE CONDITIONS MM= M+2 AND NN= N+2 ARE NOT SATISFIED,          CON 540
C EXECUTION OF THE SUBPROGRAM IS INTERRUPTED AND A MESSAGE IS     CON 550
C PRINTED OUT (AUTOMATIC ADJUSTMENT OF MM AND NN MIGHT CAUSE     CON 560
C SEVERE ERRORS IN THE MAIN PROGRAM). IF THERE IS NO DIRICH-    CON 570
C LET BOUNDARY POINT (NO C1), THE COMPUTER AUTOMATICALLY IN-    CON 580
C SERTS SUCH A POINT IN THE REGION R1 WITHIN THE BOUNDARY C2     CON 590
C (THIS MERELY PINS DOWN ONE SOLUTION OF THE NEUMANN PROBLEM).   CON 600
C IF THERE ARE MORE INTERIOR POINTS THAN ALLOWED BY THE IN-     CON 610
C PUT PARAMETER MN, EXECUTION IS INTERRUPTED AND A MESSAGE      CON 620
C IS PRINTED OUT (MN MAY NOT BE ADJUSTED AUTOMATICALLY BY      CON 630
C THE COMPUTER).                                               CON 640
      DIMENSION A(M,N), NET(MM,NN), IJ(MN,3)                   CON 650
      COMMON /SPT/ K, IJS(20,4) /CS/ NR /HV/ KK, AA(20)        CON 660
      IF (MM.NE.M+2 .OR. NN.NE.N+2) GO TO 70                   CON 670
      IP = IPRINT                                               CON 680
      IF (IP.NE.0) WRITE (6,99999)                             CON 690
C FIND PROPER STARTING POINTS FOR THE CONTOUR SCAN.              CON 700
      CALL SPOINT(M, N, MM, NN, A, NET, IP)                    CON 710
C FIND A CONTINUOUS CONTOUR PATH, IF POSSIBLE.                   CON 720
      L = 0                                                     CON 730
      10 L = L + 1                                             CON 740
      CALL CONTR(MM, NN, MN, NET, IJ, L)                       CON 750
      NP = NR                                                   CON 760
      IF (IJS(L,4).EQ.0) GO TO 50                             CON 770
      IF (L.LT.K) GO TO 10                                     CON 780
C SINCE THERE IS NO CONTINUOUS PATH FIND THE PATH WITH THE     CON 790
C LARGEST NUMBER OF BREAKS.                                     CON 800
      IG = 0                                                    CON 810
      DO 20 L=1,K                                             CON 820
        IF (IJS(L,4).GT.IG) IG = IJS(L,4)                    CON 830
      20 CONTINUE                                             CON 840
C FIND THE PATH WITH THE SMALLEST NUMBER OF BREAKS.            CON 850
      IS = IG                                                  CON 860
      DO 30 L=1,K                                             CON 870
        IF (IJS(L,4).LT.IS) IS = IJS(L,4)                    CON 880
      30 CONTINUE                                             CON 890
C CHOOSE A PATH WITH THE SMALLEST NUMBER OF BREAKS BEING      CON 900
C CLOSE TO THE HIGHEST VALUE OF DIRICHLET B.C.                 CON 910
      L = 0                                                     CON 920
      40 L = L + 1                                             CON 930
      IF (IJS(L,4).EQ.IS) GO TO 60                             CON 940
      IF (L.LT.K) GO TO 40                                     CON 950
      50 I = IJS(L,3)                                          CON 960
      IF (IP.NE.0) WRITE (6,99998) L, (IJS(L,J),J=1,2), AA(I) CON 970
      RETURN                                                  CON 980
      60 I = IJS(L,3)                                          CON 990
      IF (IP.NE.0) WRITE (6,99997) IS, AA(I), (IJS(L,J),J=1,2) CON 1000
C DETERMINE THE FINAL CONTOUR PATH.                             CON 1010
      CALL CONTR(MM, NN, MN, NET, IJ, L)                       CON 1020
      RETURN                                                  CON 1030
      70 WRITE (6,99996)                                       CON 1040
      STOP                                                    CON 1050
99999 FORMAT (1H1, 10X, 26HCONTOUR PATH OPTIMIZATION.//)      CON 1060
99998 FORMAT (/1X, 39HTHE CONTINUOUS CONTOUR SCAN HAS BEEN FO,  CON 1070
* 9HUND AFTER, 13, 15H EVALUATION(S)./17H THE STARTING POI,  CON 1080
* 8HNT IS A(, I2, 1H,, I2, 4H) AT, E12.4, 5H B.C.//)        CON 1090
99997 FORMAT (/1X, 28HTHE OPTIMUM CONTOUR SCAN HAS, I4, 2H B,  CON 1100
* 17HREAK(S) (MINIMUM)/30H AND IT STARTS FROM THE HIGHE,    CON 1110
* 17HT POSSIBLE VALUE,, E12.4/25H THE STARTING POINT IS A(,  CON 1120
* I2, 1H,, I2, 2H).//)                                       CON 1130
99996 FORMAT (1H1, 1X, 36HADJUST THE PARAMETERS MM AND NN SUCH, CON 1140
* 5H THAT/1X, 8HMM = M+2/1X, 8HNN = N+2//1X, 14HSTOP FROM CONO, CON 1150
* 3HPT.)                                                       CON 1160
      END                                                    CON 1170

      SUBROUTINE SPOINT(M, N, MM, NN, A, NET, IP)              SPO 10
C THE SUBROUTINE DETERMINES THE POSITION OF POINTS IN THE       SPO 20
C VICINITY OF THE HIGHEST AND INTERMEDIATE VALUES OF DIRICHLET SPO 30
C BOUNDARY CONDITIONS. THE ARRAYS A AND NET SERVE AS AN INPUT. SPO 40
C ML -AN INTEGER, ML=M+1                                       SPO 50
C NL -AN INTEGER, NL=N+1                                       SPO 60
C THE PARAMETERS ARE DEFINED IN CONOPT.                          SPO 70
      DIMENSION A(M,N), NET(MM,NN)                             SPO 80
      COMMON /COEF/ ML, NL /SPT/ K, IJS(20,4) /HV/ KK, AA(20) SPO 90
C FIND HIGH VALUE DIRICHLET B.C.                                SPO 100
      CALL HIVA(M, N, MM, NN, A, NET, IP)                      SPO 110
      K = 0                                                     SPO 120
      DO 100 L=1, KK                                           SPO 130

```

```

      AM = AA(L)
C FIND POINTS SEPARATED FROM THE HIGH VALUE DIRICHLET B.C.
C BY ONE MESH LENGTH. THESE POINTS LIE ON THE ROWS.
      DO 30 I=1,M
        DO 20 J=1,N
          IF (NET(I+1,J+1).NE.0) GO TO 20
          IF (A(I-1,J).EQ.AM .AND. NET(I,J).NE.1) GO TO 10
          IF (A(I+1,J).NE.AM .OR. NET(I+2,J+2).EQ.1) GO TO 20
10      K = K + 1
          IF (K.GT.20) GO TO 110
          IJS(K,1) = I
          IJS(K,2) = J
          IJS(K,3) = L
20      CONTINUE
30      CONTINUE
C FIND POINTS AS ABOVE BUT LYING ON THE COLUMNS.
      DO 60 J=1,N
        DO 50 I=1,M
          IF (NET(I+1,J+1).NE.0) GO TO 50
          IF (A(I,J+1).EQ.AM .AND. NET(I,J+2).NE.1) GO TO 40
          IF (A(I,J-1).NE.AM .OR. NET(I+2,J).EQ.1) GO TO 50
40      K = K + 1
          IF (K.GT.20) GO TO 110
          IJS(K,1) = I
          IJS(K,2) = J
          IJS(K,3) = L
50      CONTINUE
60      CONTINUE
C FIND POINTS ADJACENT TO DIRICHLET B.C., LOCATED AT CORNERS
      DO 90 I=2,ML
        DO 80 J=2,NL
          IF (NET(I,J).NE.0) GO TO 80
          IF (A(I-2,J-1).EQ.AM .AND. NET(I-1,J-1).EQ.1 .AND.
*          NET(I,J-1).EQ.1) GO TO 70
          IF (A(I-1,J).EQ.AM .AND. NET(I-1,J+1).EQ.1 .AND.
*          NET(I-1,J).EQ.1) GO TO 70
          IF (A(I,J-1).EQ.AM .AND. NET(I+1,J+1).EQ.1 .AND.
*          NET(I,J+1).EQ.1) GO TO 70
          IF (A(I-1,J-2).NE.AM .OR. NET(I+1,J-1).NE.1 .OR.
*          NET(I+1,J).NE.1) GO TO 80
70      K = K + 1
          IF (K.GT.20) GO TO 110
          IJS(K,1) = I - 1
          IJS(K,2) = J - 1
          IJS(K,3) = L
80      CONTINUE
90      CONTINUE
100 CONTINUE
110 IF(IP.NE.0) WRITE(6,99999) (I,(IJS(I,J), J=1,2), I=1,K)
      RETURN
99999 FORMAT (/IX, 39HSTARTING POINTS FOR CONTOUR OPTIMIZATIO, 1HN//
* 5X, 2HNO, 5X, 1HI, 5X, 1HJ/(I7, 2I6))
      END

```

```

      SUBROUTINE HIVA(M, N, MM, NN, A, NET, IP)
C THE SUBROUTINE DETERMINES THE HIGH VALUE DIRICHLET BOUNDARY
C CONDITIONS AND IT RETURNS THEM IN THE ARRAY AA.
C PARAMETERS ARE DEFINED IN SPOINT.
      DIMENSION A(M,N), NET(MM,NN)
      COMMON /HV/ KK, AA(20)
C FIND THE MAXIMUM AND MINIMUM VALUES IN THE ARRAY A.
C ARE THERE ANY POINTS OF TYPE S1
      KK = 0
      DO 20 J=1,N
        J1 = J + 1
        DO 10 I=1,M
          IF (NET(I+1,J1).EQ.1) KK = KK + 1
10      CONTINUE
20      CONTINUE
C IF KK=0 (NO S1 POINTS), THE SOLUTION WILL BE PINNED DOWN.
      IF (KK) 30, 30, 60
30      DO 50 J=1,N

```


J1 = J + 1	HIV 190
DO 40 I=1,M	HIV 200
IF (NET(I+1,J1).NE.0) GO TO 40	HIV 210
NET(I+1,J1) = 1	HIV 220
A(I,J) = 50.0	HIV 230
AA(1) = 50.0	HIV 240
KK = 1	HIV 250
GO TO 120	HIV 260
40 CONTINUE	HIV 270
50 CONTINUE	HIV 280
WRITE (6,99999)	HIV 290
STOP	HIV 300
60 AM = 0.0	HIV 310
AS = 100.0	HIV 320
DO 80 J=1,N	HIV 330
J1 = J + 1	HIV 340
DO 70 I=1,M	HIV 350
IF (NET(I+1,J1).NE.1) GO TO 70	HIV 360
IF (A(I,J).GT.AM) AM = A(I,J)	HIV 370
IF (A(I,J).LT.AS) AS = A(I,J)	HIV 380
70 CONTINUE	HIV 390
80 CONTINUE	HIV 400
KK = 1	HIV 410
AA(1) = AM	HIV 420
C FIND INTERMEDIATE VALUES.	HIV 430
DO 110 K=1,19	HIV 440
AM = AS	HIV 450
DO 100 J=1,N	HIV 460
J1 = J + 1	HIV 470
DO 90 I=1,M	HIV 480
IF (NET(I+1,J1).NE.1) GO TO 90	HIV 490
IF (A(I,J).GT.AM .AND. A(I,J).LT.AA(K)) AM = A(I,J)	HIV 500
90 CONTINUE	HIV 510
100 CONTINUE	HIV 520
IF (AM.EQ.AS) GO TO 120	HIV 530
KK = KK + 1	HIV 540
AA(KK) = AM	HIV 550
110 CONTINUE	HIV 560
120 IF (IP.NE.0) WRITE (6,99998) (I,AA(I),I=1,KK)	HIV 570
RETURN	HIV 580
99999 FORMAT (/1X, 37HTHERE ARE NO FREE POINTS WITHIN A AND,	HIV 590
* 26H NET. SOLUTION IMPOSSIBLE.//1X, 14HSTOP FROM HIVA,	HIV 600
* 11H IN CONOPT.)	HIV 610
99998 FORMAT (/1X, 25HHIGH VALUE DIRICHLET B.C./(I7, E15.5))	HIV 620
END	HIV 630
SUBROUTINE CONTR(MM, NN, MN, NET, IJ, L1)	CON 10
C THE SUBROUTINE CONTR TRACES A CONTOUR SCAN WHICH COVERS	CON 20
C ONLY THE INTERIOR OF THE REGION DEFINED BY BOUNDARY OF AN	CON 30
C ARBITRARY SHAPE. THE SUBROUTINE REQUIRES A STARTING POINT,	CON 40
C I,J,AND IT RETURNS THE CONTOUR PATH STORED IN THE ARRAY IJ	CON 50
C IX -AN ARRAY USED ONLY IN THIS SUBROUTINE FOR SETTING	CON 60
C A DIRECTIONAL PATTERN OF I.	CON 70
C JY -AS ABOVE, BUT FOR J.	CON 80
C L -AN INDICATOR OF THE PERFORMED TURN. L=1-TURN LEFT,	CON 90
C L=2-GO STRAIGHT, L=3-TURN RIGHT.	CON 100
C K -AN INDICATOR OF THE CHOSEN DIRECTION.K IS RELATED TO L	CON 110
C OTHER PARAMETERS AS DEFINED IN CONOPT.	CON 120
DIMENSION NET(MM,NN), IJ(MN,3)	CON 130
COMMON /COEF/ ML, NL /TEMP/ IX(5), JY(5)	CON 140
COMMON /SPT/ K1, IJS(20,4) /CS/ NR	CON 150
C SET INITIAL PARAMETERS.	CON 160
I = IJS(L1,1) + 1	CON 170
J = IJS(L1,2) + 1	CON 180
L = 2	CON 190
NR = 0	CON 200
IJS(L1,4) = 0	CON 210
GO TO 30	CON 220
C FIND INITIAL DIRECTION OF THE SCAN.	CON 230
10 IX(5) = IX(1)	CON 240
JY(5) = JY(1)	CON 250
DO 20 K=1,4	CON 260
I = IX(K+1)	CON 270
J = JY(K+1)	CON 280

```

        I1 = IX(K)
        J1 = JY(K)
        IF (NET(I,J).NE.1 .AND. NET(I1,J1).EQ.1) GO TO 30
20 CONTINUE
30 NR = NR + 1
    IF (NR.GT.MN) GO TO 110
C SAVE I,J, AND NET(I,J). PROTECT NET(I,J).
    IJ(NR,1) = I - 1
    IJ(NR,2) = J - 1
    IJ(NR,3) = NET(I,J)
    NET(I,J) = 1
C SET THE DIRECTION PATTERN.
    IX(1) = I - 1
    JY(1) = J
    IX(2) = I
    JY(2) = J + 1
    IX(3) = I + 1
    JY(3) = J
    IX(4) = I
    JY(4) = J - 1
    IF (NR.EQ.1) GO TO 10
C COMPUTE THE RELATIVE DIRECTION.
    K2 = K + L - 2
    K = K2 + 4*(1-((3+K2)/4))
    K2 = K - 1
    IF (K2.EQ.0) GO TO 60
C ROTATE THE PATTERN ACCORDING TO THE RELATIVE DIRECTION.
    DO 50 L2=1,K2
        IX(5) = IX(1)
        JY(5) = JY(1)
    DO 40 L=1,4
        IX(L) = IX(L+1)
        JY(L) = JY(L+1)
40 CONTINUE
50 CONTINUE
C CHECK ALL THE POSSIBLE MOVES. IF A MOVE IS POSSIBLE, THE
C VALUE OF L WILL BE PRESERVED.
60 DO 70 L=1,3
    I = IX(L)
    J = JY(L)
    IF (NET(I,J).NE.1) GO TO 30
70 CONTINUE
C FIND AND COUNT POSSIBLE BREAKS IN THE SCAN.
    K = 1
    L = 2
    DO 90 I=2,ML
        DO 80 J=2,NL
            IF (NET(I,J).NE.1) IJS(L1,4) = IJS(L1,4) + 1
            IF (NET(I,J).NE.1) GO TO 30
80 CONTINUE
90 CONTINUE
C RESET THE ARRAY NET
    DO 100 K=1,NR
        I = IJ(K,1) + 1
        J = IJ(K,2) + 1
        NET(I,J) = IJ(K,3)
100 CONTINUE
    RETURN
110 WRITE (6,99999)
    STOP
99999 FORMAT (/1X, 35HTHERE ARE MORE INTERIOR POINTS THAN, 6H ALLOW,
* 9HED BY MN./1X, 19HINCREASE MN TO M*N./1X, 14HSTOP FROM CONT,
* 12HR IN CONOPT.)
    END

```

```

SUBROUTINE RASTER(MM, NN, MN, NET, IJ)
C THE SUBROUTINE RETURNS A RASTER SCAN OF THE INTERIOR OF AN MXN
C ARRAY. ONLY ACTIVE POINTS OF THE SCAN ARE SAVED IN IJ.
    DIMENSION NET(MM,NN), IJ(MN,3)
    COMMON /COEF/ ML, NL /CS/ NR
    WRITE (6,99999)
    NR = 0
    DO 20 I=2,ML
        RAS 10
        RAS 20
        RAS 30
        RAS 40
        RAS 50
        RAS 60
        RAS 70
        RAS 80

```

```

          DO 10 J=2,NL
            IF (NET(I,J).EQ.1) GO TO 10
            NR = NR + 1
            IJ(NR,1) = I - 1
            IJ(NR,2) = J - 1
10      CONTINUE
20      CONTINUE
        RETURN
99999 FORMAT (1H1, 10X, 11HRASTER SCAN)
        END

```

```

RAS 90
RAS 100
RAS 110
RAS 120
RAS 130
RAS 140
RAS 150
RAS 160
RAS 170
RAS 180

```

```

          SUBROUTINE SPIRAL(M, N, MM, NN, MN, NET, IJ)
C THE SUBROUTINE TRACES AN INWARD SPIRAL PATH WITHIN AN
C M X N ARRAY A. THE SCAN IS RETURNED IN IJ.
C THE PARAMETERS ARE DEFINED IN THE SUBROUTINE CONOPT AND SPOINT.
          DIMENSION NET(MM,NN), IJ(MN,3)
          COMMON /COEF/ ML, NL /CS/ NR
          WRITE (6,99999)
          NR = 0
C FIND THE NUMBER OF RECTANGLES WHICH WILL CONSTITUTE
C THE SPIRAL SCAN.
          IH = MIN0(ML,NL)/2
C BEGIN FROM THE LARGEST RECTANGLE AND SUCCESSIVELY REDUCE
C ITS SIZE.
          DO 90 I1=1,IH
            IS = I1
            K = 0
C LOCATE START AND END POINTS FOR EACH SIDE OF A RECTANGLE.
          DO 80 I2=1,2
            I = IS
            IF (I2.EQ.2) I = IE + 1
            IE = N - IS
            DO 70 I3=1,2
              K = K + 1
C SCAN ONE SIDE OF A RECTANGLE AND DISTRIBUTE THE SEQUENCE OF
C POINTS TO I OR TO J.
          DO 60 I4=IS,IE
            GO TO (10, 20, 30, 40), K
10          J = I4
            GO TO 50
20          I = I4
            GO TO 50
30          J = NL - I4
            GO TO 50
40          I = ML - I4
50          IF (NET(I+1,J+1).EQ.1) GO TO 60
            NR = NR + 1
            IJ(NR,1) = I
            IJ(NR,2) = J
60          CONTINUE
            J = IE + 1
            IF (I2.EQ.2) J = IS
            IE = M - IS
70          CONTINUE
80          CONTINUE
90          CONTINUE
        RETURN
99999 FORMAT (1H1, 10X, 11HSPIRAL SCAN)
        END

```

```

SPI 10
SPI 20
SPI 30
SPI 40
SPI 50
SPI 60
SPI 70
SPI 80
SPI 90
SPI 100
SPI 110
SPI 120
SPI 130
SPI 140
SPI 150
SPI 160
SPI 170
SPI 180
SPI 190
SPI 200
SPI 210
SPI 220
SPI 230
SPI 240
SPI 250
SPI 260
SPI 270
SPI 280
SPI 290
SPI 300
SPI 310
SPI 320
SPI 330
SPI 340
SPI 350
SPI 360
SPI 370
SPI 380
SPI 390
SPI 400
SPI 410
SPI 420
SPI 430
SPI 440
SPI 450
SPI 460
SPI 470
SPI 480

```

```

          SUBROUTINE SORCAR(M, N, MN, A, IJ, C, IP, ER)
C THIS SUBROUTINE OPTIMIZES THE OVERRELAXATION FACTOR USING CARRE*S
C TECHNIQUE.
          DIMENSION A(M,N), IJ(MN,3), C(MN), GX(3)
          COMMON /LOG/ Q, QQ
          LOGICAL Q, QQ, QQQ
          WRITE (6,99999)
C THE FIRST ITERATION USES BETA=1.0
          BETA = 1.
C THE FIRST INTERVAL OF INT ITERATIONS USES BETA=1.375
          BOPT = 1.375
C SET THE INTERVAL OF ITERATIONS

```

```

SOR 10
SOR 20
SOR 30
SOR 40
SOR 50
SOR 60
SOR 70
SOR 80
SOR 90
SOR 100
SOR 110
SOR 120

```

```

      INT = 12
      Q = .TRUE.
      QQQ = .FALSE.
      IN = 0
      SUM = 0.0
      JOB = 0
10  JOB = JOB + 1
      IN = IN + 1
      QQ = IN.LT.INT
      BIG = 0.0
      S = SUM
      SUM = 0.0
      CALL SOR(M, N, MN, A, LJ, C, BETA, SUM, BIG)
      QQ = (.NOT.QQ) .AND. QQQ
C Q=.FALSE. IF BETA OPTIMUM IS FOUND.
      IF (.NOT.Q) GO TO 60
C QQ=.TRUE. IF JOB.GT.INT+1 AND IN=INT
      IF (QQ) GO TO 30
C QQQ=.TRUE. IF JOB.GT.INT+1
      IF (QQQ) GO TO 10
      IF (JOB.EQ.1) BETA = BOPT
      IF (JOB.EQ.1) IN = 0
      IF (IN.LE.INT-3) GO TO 10
C COMPUTE THREE SUCCESSIVE RATIOS OF THE DISPLACEMENT VECTOR NORMS
C AT THE END OF THE FIRST INTERVAL
      L = IN - (INT-3)
      GX(L) = SUM/S
      IF (IN.LT.INT) GO TO 10
      QQQ = .TRUE.
      D1 = GX(1) - GX(2)
      D2 = GX(2) - GX(3)
C CRITERION FOR DETERMINING THE FEASIBILITY OF AITKEN*S
C EXTRAPOLATION, THE DIFFERENCES BETWEEN THE RATIOS MUST DECREASE
C AND HAVE THE SAME SIG
      IF (ABS(D1).LE.ABS(D2) .OR. (D1/D2).LT.0.0) GO TO 20
      GAM = GX(1) - ((GX(2)-GX(1))**2/(GX(1)-2.*GX(2)+GX(3)))
      GO TO 40
20  WRITE (6,99998)
30  GAM = SUM/S
40  IN = 0
      BL = BOPT
C COMPUTE A NEW ESTIMATE OF BETA OPTIMUM
      BOPT = 2./(1.+SQRT(ABS(1.-(GAM+BETA-1.)**2/(GAM*BETA**2))))
C THE FOLLOWING BOPT REDUCTION PREVENTS OVERESTIMATION OF BETA
      BETA = BOPT - (2.-BOPT)/4.
C CRITERION FOR STOPPING THE PROCESS OF IMPROVING BETA.
      IF (ABS(BOPT-BL)/(2.-BOPT).GT.0.05) GO TO 50
C BETA OPTIMUM HAS BEEN FOUND
      Q = .FALSE.
      IN = INT
      BETA = BOPT
      GAM = BETA - 1.
      GR = GAM/(2.-BETA)
      WRITE (6,99997) BETA, JOB
50  IF (GAM.EQ.1.) GO TO 10
      ERM = BIG*GAM/(1.-GAM)
      GO TO 70
60  ERM = BIG*GR
70  IF (ABS(ERM).LE.ER) GO TO 80
      IF (JOB.LT.IP) GO TO 10
80  WRITE (6,99996) JOB, ERM
C IF BETA OPTIMUM HAS NOT BEEN FOUND YET, BUT ERM.LE.ER, PRINT BETA.
      IF (Q) WRITE (6,99995) BETA
      RETURN
99999 FORMAT (/1X, 30HSOLUTION BY SOR (CARRE) METHOD)
99998 FORMAT (/5X, 45HAITKEN*S EXTRAPOLATION NOT FEASIBLE FOR BETA ,
* 13HOPTIMIZATION.)
99997 FORMAT (5X, 14HBETA OPTIMUM =, E14.7, 6H AFTER, I4, 7H ITERAT,
* 5HIIONS.)
99996 FORMAT (/10X, 22HNUMBER OF ITERATIONS =, I4/10X, 9HMAXIMUM E,
* 4HRROR, 8X, 1H=, E10.3, 8H PERCENT)
99995 FORMAT (10X, 22HBETA (NOT OPTIMUM) =, E14.7)
      END

```

SOR 130
 SOR 140
 SOR 150
 SOR 160
 SOR 170
 SOR 180
 SOR 190
 SOR 200
 SOR 210
 SOR 220
 SOR 230
 SOR 240
 SOR 250
 SOR 260
 SOR 270
 SOR 280
 SOR 290
 SOR 300
 SOR 310
 SOR 320
 SOR 330
 SOR 340
 SOR 350
 SOR 360
 SOR 370
 SOR 380
 SOR 390
 SOR 400
 SOR 410
 SOR 420
 SOR 430
 SOR 440
 SOR 450
 SOR 460
 SOR 470
 SOR 480
 SOR 490
 SOR 500
 SOR 510
 SOR 520
 SOR 530
 SOR 540
 SOR 550
 SOR 560
 SOR 570
 SOR 580
 SOR 590
 SOR 600
 SOR 610
 SOR 620
 SOR 630
 SOR 640
 SOR 650
 SOR 660
 SOR 670
 SOR 680
 SOR 690
 SOR 700
 SOR 710
 SOR 720
 SOR 730
 SOR 740
 SOR 750
 SOR 760
 SOR 770
 SOR 780
 SOR 790
 SOR 800
 SOR 810
 SOR 820
 SOR 830
 SOR 840
 SOR 850

SUBROUTINE SOR(M, N, MN, A, IJ, C, BETA, SUM, BIG)	SOR	10
C THIS IS THE SUCCESSIVE OVERRELAXATION ALGORITHM.	SOR	20
C THE SUBROUTINE PERFORMS ONE COMPLETE SWEEP OF THE INTERIOR, AND	SOR	30
C COMPUTES THE ERROR CRITERION COMPONENTS.	SOR	40
DIMENSION A(M,N), IJ(MN,3), C(MN)	SOR	50
COMMON /CS/ NR /LOG/ Q, QQ	SOR	60
EXTERNAL AV	SOR	70
LOGICAL Q, QQ	SOR	80
DO 10 K=1,NR	SOR	90
I = IJ(K,1)	SOR	100
J = IJ(K,2)	SOR	110
L = IJ(K,3)	SOR	120
AL = A(I,J)	SOR	130
AN = AV(M,N,MN,A,C,I,J,K,L)	SOR	140
AN = AL + BETA*(AN-AL)	SOR	150
A(I,J) = AN	SOR	160
RES = ABS(AN-AL)	SOR	170
IF (Q) SUM = SUM + RES	SOR	180
IF (QQ) GO TO 10	SOR	190
IF (RES.GT.BIG) BIG = RES	SOR	200
10 CONTINUE	SOR	210
RETURN	SOR	220
END	SOR	230
FUNCTION AV(M, N, MN, A, C, I, J, K, L)	AV	10
DIMENSION A(M,N), C(MN)	AV	20
GO TO (10, 20, 30, 40, 50), L	AV	30
C NEUMANN B.C. ARE PARALLEL TO J-AXIS	AV	40
10 AV = (2.*A(I+1,J)+C(K)+C(K)+A(I,J+1)+A(I,J-1))/4.	AV	50
RETURN	AV	60
C NEUMANN B.C. ARE PARALLEL TO I-AXIS	AV	70
20 AV = (2.*A(I,J+1)+C(K)+C(K)+A(I+1,J)+A(I-1,J))/4.	AV	80
RETURN	AV	90
C NEUMANN B.C. AS IN 1, BUT FOR I.GT.IN 1	AV	100
30 AV = (2.*A(I-1,J)+C(K)+C(K)+A(I,J+1)+A(I,J-1))/4.	AV	110
RETURN	AV	120
C NEUMANN B.C. AS IN 2, BUT FOR J.GT. IN 2	AV	130
40 AV = (2.*A(I,J-1)+C(K)+C(K)+A(I+1,J)+A(I-1,J))/4.	AV	140
RETURN	AV	150
C ORDINARY 5-POINT OPERATOR FOR THE INTERIOR.	AV	160
50 AV = (A(I,J+1)+A(I+1,J)+A(I-1,J)+A(I,J-1))/4.	AV	170
RETURN	AV	180
END	AV	190
SUBROUTINE BOUND(M, N, MM, NN, A, NET)	BOU	10
C THE SUBROUTINE READS THE BOUNDARY CONDITIONS SPECIFIED IN THE	BOU	20
C DATA CARDS AND SETS THEM INTO THE ARRAYS A AND NET. THE BOUNDARY	BOU	30
C CONDITIONS ARE DETERMINED BY SEGMENTS P1P2 WHERE POINTS P1(I1,J1)	BOU	40
C AND P2(I2,J2) COINCIDE WITH THE MESH POINTS IMPOSED ON THE REGION	BOU	50
C A(M,N). THE POINTS MUST SATISFY THE FOLLOWING CONDITION P1.LE.P2.	BOU	60
C THE VALUES AB ASSOCIATED WITH THE POINTS P1 AND P2 CAN BE AB1.LE.AB2	BOU	70
C OR AB1.GE.AB2. DIRICHLET B.C. ARE DISTINGUISHED FROM NEUMANN B.C.	BOU	80
C BY THE INDICATOR IND WHERE IND=1 FOR DIRICHLET OR IND=2 FOR NEUMANN	BOU	90
C B.C. EACH DATA CARD CONTAINS ONLY ONE SEGMENT P1P2 (OR A POINT).	BOU	100
C THE INDEX NEXT.NE.0 TERMINATES THE READING PROCESS FOR ONE PROBLEM.	BOU	110
DIMENSION A(M,N), NET(MM,NN)	BOU	120
COMMON /COEF/ ML, NL /TEMP/ ID(5), IY(5)	BOU	130
C CLEAR THE ARRAYS.	BOU	140
DO 20 J=1,N	BOU	150
DO 10 I=1,M	BOU	160
A(I,J) = 0.0	BOU	170
NET(I+1,J+1) = 0	BOU	180
10 CONTINUE	BOU	190
20 CONTINUE	BOU	200
C SET BORDERS IN THE ARRAYS A AND NET.	BOU	210
DO 30 I=2,ML	BOU	220
NET(I,2) = 1	BOU	230
NET(I,NL) = 1	BOU	240
30 CONTINUE	BOU	250
DO 40 J=2,NL	BOU	260
NET(2,J) = 1	BOU	270

```

      NET(ML,J) = 1
40 CONTINUE
      DO 50 I=1,MM
          NET(I,1) = 1
          NET(I,NN) = 1
50 CONTINUE
      DO 60 J=1,NN
          NET(1,J) = 1
          NET(MM,J) = 1
60 CONTINUE
70 READ (5,99999) I1, J1, AB1, I2, J2, AB2, IND, NEXT
      IF (I1.GT.I2 .OR. J1.GT.J2) GO TO 150
C PLACE BOUNDARY PARALLEL TO I-AXIS INTO A.
      IF (I1.EQ.I2) GO TO 90
      AB = (AB2-AB1)/FLOAT(I2-I1)
      DO 80 I=I1,I2
          A(I,J1) = AB1 + AB*FLOAT(I-I1)
C THE FOLLOWING STATEMENT TRUNCATES ABOVE 5TH DECIMAL DIGIT.
      IF (AB.NE.0.0) A(I,J1) = AINT(A(I,J1)*1.E5)*1.E-5
      NET(I+1,J1+1) = IND
80 CONTINUE
      GO TO 120
C PLACE BOUNDARY PARALLEL TO J-AXIS INTO A.
90 IF (J1.EQ.J2) GO TO 110
      AB = (AB2-AB1)/FLOAT(J2-J1)
      DO 100 J=J1,J2
          A(I1,J) = AB1 + AB*FLOAT(J-J1)
          IF (AB.NE.0.0) A(I1,J) = AINT(A(I1,J)*1.E5)*1.E-5
          NET(I1+1,J+1) = IND
100 CONTINUE
      GO TO 120
C PLACE THE POINT (I1,J1) INTO A.
110 A(I1,J1) = AB1
      NET(I1+1,J1+1) = IND
120 IF (NEXT.EQ.0) GO TO 70
C PRINT THE ARRAY NET.
      WRITE (6,99998)
      ND = 1 + (N-1)/10
      DO 130 J=1,ND
          ID(J) = J
130 CONTINUE
      WRITE (6,99997) (ID(J),J=1,ND)
      DO 140 I=1,M
          WRITE (6,99996) I, (NET(I+1,J),J=2,NL)
140 CONTINUE
      RETURN
150 WRITE (6,99995)
      STOP
99999 FORMAT (2(2I5, F10.4), 2I5)
99998 FORMAT (1H1, 10X, 40HTHE FOLLOWING MATRIX REPRESENTS THE NET,,
      * 5X, 11H0=VARIABLES/56X, 26H1=DIRICHLET BOUNDARY COND./56X,
      * 24H2=NEUMANN BOUNDARY COND./)
99997 FORMAT (1H , 35X, 5(I1, 19X)/)
99996 FORMAT (11X, I2, 4X, 50I2)
99995 FORMAT (1H1, 1X, 41HINCORRECT INPUT DATA, I1.GT.I2 OR J1.GT.J,
      * 2H2.)
      END

```

```

      SUBROUTINE DINE(M, N, MM, NN, MN, A, NET, IJ, C)
C THE SUBROUTINE RETURNS REFERENCE MARKS LOCATED IN THE ARRAY IJ(I,3).
C THE MARKS WILL ACCOMPANY THE SCAN AND HENCE SIMPLIFY THE LOGIC OF
C THE ITERATION PROCESS. BOUNDARY CONDITIONS IN THE ARRAY NET ARE
C IDENTIFIED BY 1= DIRICHLET, 2= NEUMANN, AND 0= INTERIOR OF THE
C REGION.
      DIMENSION A(M,N), NET(MM,NN), IJ(MN,3), C(MN)
      COMMON /CS/ NR
C SET MARKS OF REFERENCE, AND SAVE NEUMANN B.C. FROM A INTO C.
C 1-4 = NEUMANN B.C. LOCATED AT (I-1,J), (I,J-1), (I+1,J), AND (I,J+1)
C WITH RESPECT TO THE INTERIOR OF THE REGION, RESPECTIVELY.
C 5-INTERNAL (REGULAR) POINTS. NOTE THAT THE MARK CHANGE FROM 0 TO 5
C IS INTERNAL TO THE PROGRAM AND IS INTRODUCED TO SIMPLIFY THE
C SELECTION OF EQUATIONS PERTINENT TO EITHER THE NEUMANN POINTS (1-4)
C OR THE REGULAR POINTS (5) /SEE FUNCTION AV/.

```

```

DO 30 K=1, NR
  I = IJ(K,1) + 1
  J = IJ(K,2) + 1
  IF (NET(I,J)-1) 10, 40, 20
10 IJ(K,3) = 5
  C(K) = 0.0
  GO TO 30
20 IF (NET(I+1,J).EQ.0) IJ(K,3) = 1
  IF (NET(I,J+1).EQ.0) IJ(K,3) = 2
  IF (NET(I-1,J).EQ.0) IJ(K,3) = 3
  IF (NET(I,J-1).EQ.0) IJ(K,3) = 4
  C(K) = A(I-1,J-1)
30 CONTINUE
  WRITE (6,99999) NR
  RETURN
40 WRITE (6,99998) I, J
  STOP
99999 FORMAT (/10X, 30HNUMBER OF SCANNED POINTS, NR= , I5)
99998 FORMAT (/1X, 36HINCORRECT SCANNING SEQUENCE AT NET (, I2, 1H,,
* I2, 2H).)
END

```

DIN 160
DIN 170
DIN 180
DIN 190
DIN 200
DIN 210
DIN 220
DIN 230
DIN 240
DIN 250
DIN 260
DIN 270
DIN 280
DIN 290
DIN 300
DIN 310
DIN 320
DIN 330
DIN 340
DIN 350
DIN 360

```

SUBROUTINE INVA(M, N, MM, NN, A, NET)
C THE SUBROUTINE READS A CONSTANT INITIAL VALUE, AND INTRODUCES IT
C INTO THE ARRAY A. THE INITIAL VALUE DATUM FOLLOWS THE BOUNDARY
C CONDITION DATA.
  DIMENSION A(M,N), NET(MM,NN)
  READ (5,99999) V
  DO 20 J=1,N
    DO 10 I=1,M
      IF (NET(I+1,J+1).NE.1) A(I,J) = V
10 CONTINUE
20 CONTINUE
  WRITE (6,99998) V
  RETURN
99999 FORMAT (F10.0)
99998 FORMAT (10X, 29HCONSTANT INITIAL VALUE, V=, F7.2)
END

```

INV 10
INV 20
INV 30
INV 40
INV 50
INV 60
INV 70
INV 80
INV 90
INV 100
INV 110
INV 120
INV 130
INV 140
INV 150
INV 160

```

SUBROUTINE PRES(M, N, A)
C THE SUBROUTINE PRINTS NUMERICAL RESULTS FROM THE ARRAY A.
C EACH PRINTED PAGE WILL CONTAIN 10 COLUMNS OF A.
  DIMENSION A(M,N)
  COMMON /TEMP/ LA(10)
  DO 30 L=1,N,10
    JS = L - 1
    JE = L + 9
    DO 10 I=1,10
      LA(I) = JS + I
10 CONTINUE
  WRITE (6,99999) LA
  DO 20 I=1,M
    WRITE (6,99998) I, (A(I,J),J=L,JE)
20 CONTINUE
30 CONTINUE
  RETURN
99999 FORMAT (1H1, 15X, 10(I3, 9X)/)
99998 FORMAT (1H , 5H A(, I2, 1H,, 3X, 10(F10.5, 2X))
END

```

PRE 10
PRE 20
PRE 30
PRE 40
PRE 50
PRE 60
PRE 70
PRE 80
PRE 90
PRE 100
PRE 110
PRE 120
PRE 130
PRE 140
PRE 150
PRE 160
PRE 170
PRE 180
PRE 190
PRE 200

```

SUBROUTINE MAP(M, N, MM, NN, A, NET)
C THE SUBROUTINE PRINTS THE EQUIPOTENTIAL MAP.
  DIMENSION A(M,N), NET(MM,NN), CA(10), SP(10)
  COMMON /TEMP/ IZ(10)
  DATA CA(1), CA(2), CA(3), CA(4), CA(5), CA(6), CA(7), CA(8),
* CA(9), CA(10) /1H0,1H1,1H2,1H3,1H4,1H5,1H6,1H7,1H8,1H9/,

```

MAP 10
MAP 20
MAP 30
MAP 40
MAP 50
MAP 60

* ZERO, VAL, BLANK /1HZ,1H\$,1H /	MAP	70
C READ EQUIPOTENTIALS TO BE PRINTED OUT.	MAP	80
READ (5,99999) (SP(I),I=1,10), TOL	MAP	90
C PREPARE THE MAP.	MAP	100
DO 50 J=1,N	MAP	110
DO 40 I=1,M	MAP	120
IF (NET(I+1,J+1).NE.1) GO TO 20	MAP	130
IF (A(I,J).EQ.0.0) GO TO 10	MAP	140
A(I,J) = VAL	MAP	150
GO TO 40	MAP	160
10 A(I,J) = ZERO	MAP	170
GO TO 40	MAP	180
20 DO 30 K=1,10	MAP	190
IF (A(I,J).LT.SP(K)-TOL .OR. A(I,J).GT.SP(K)+TOL) GO TO	MAP	200
*	MAP	210
30	MAP	220
A(I,J) = CA(K)	MAP	230
GO TO 40	MAP	240
30 CONTINUE	MAP	250
A(I,J) = BLANK	MAP	260
40 CONTINUE	MAP	270
50 CONTINUE	MAP	280
C WRITE THE MAP.	MAP	290
WRITE (6,99998)	MAP	300
DO 60 I=1,M	MAP	310
WRITE (6,99997) (A(I,J),J=1,N)	MAP	320
60 CONTINUE	MAP	330
C WRITE THE LIMITS FOR THE EQUIPOTENTIAL STRIPS IN THE PRINTED MAP.	MAP	340
WRITE (6,99996)	MAP	350
DO 70 I=1,10	MAP	360
IZ(I) = I - 1	MAP	370
70 CONTINUE	MAP	380
DO 80 I=1,5	MAP	390
P1 = SP(I) - TOL	MAP	400
P2 = SP(I) + TOL	MAP	410
P3 = SP(I+5) - TOL	MAP	420
P4 = SP(I+5) + TOL	MAP	430
IF (P1.LT.0.) P1 = 0.	MAP	440
WRITE (6,99995) IZ(I), P1, P2, IZ(I+5), P3, P4	MAP	450
80 CONTINUE	MAP	460
WRITE (6,99994)	MAP	470
RETURN	MAP	480
C IF THE ARRAY A IS LARGER THAN 30X30, CHANGE THE FORMAT 99997	MAP	490
C TO (15X,50(A1,1X)).	MAP	500
99999 FORMAT (10F7.3, F10.3)	MAP	510
99998 FORMAT (1H1, 7X, 41H THE FOLLOWING IS A CONTOUR MAP OF THE EQU,	MAP	520
* 31HIPOTENTIALS WITHIN GIVEN LIMITS/)	MAP	530
99997 FORMAT (15X, 20(A1, 2X)/)	MAP	540
99996 FORMAT (/13X, 8HCONTOURS)	MAP	550
99995 FORMAT (13X, I1, 1H=, E11.4, 4H TO , E11.4, 6H ** , I1, 1H=,	MAP	560
* E11.4, 4H TO , E11.4)	MAP	570
99994 FORMAT (13X, 10HBOUNDARIES/13X, 22HZ= ZERO DIRICHLET B.C.,	MAP	580
* 8X, 4H** , 25H\$= NONZERO DIRICHLET B.C.)	MAP	590
END		

ALGORITHM 500

Minimization of Unconstrained Multivariate Functions [E4]

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Key Words and Phrases: minimization, optimization

CR Categories: 5.15, 5.19

Language: Fortran

DESCRIPTION

Purpose

This subroutine finds a local minimum of a nonlinear function of n variables $f(x)$ where $x = (x_1, x_2, \dots, x_n)$, $n > 1$, can be any real numbers.

Methods

Quasi-Newton methods are iterative methods which approximate Newton's method without calculating second derivatives. The iterative sequence is defined by

$$x^{(k+1)} = x^{(k)} - \alpha^{(k)} H^{(k)} g^{(k)},$$

where $g^{(k)} = \nabla f(x^{(k)})$, the gradient of f at $x^{(k)}$, $H^{(k)}$ is a matrix which is designed to approximate the inverse Hessian matrix of f at $x^{(k)}$, and $\alpha^{(k)}$ is an appropriately chosen scalar. The sequence of matrices $H^{(k)}$ is chosen to satisfy the quasi-Newton equation $H^{(k+1)} y^{(k)} = \sigma^{(k)}$, where $\sigma^{(k)} = x^{(k+1)} - x^{(k)}$ and $y^{(k)} = g^{(k+1)} - g^{(k)}$. In general, $H^{(k+1)}$ is generated by $H^{(k+1)} = H^{(k)} + D^{(k)}$, where $D^{(k)}$ is chosen to satisfy the equation $D^{(k)} y^{(k)} = \sigma^{(k)} - H^{(k)} y^{(k)}$.

The choice of $\alpha^{(k)}$ can be accomplished either by a linear-search or a step-length method. Perhaps the most complete current reference to the general theory is Powell [8]. Two different methods for determining $\alpha^{(k)}$ are incorporated and are outlined below. The Broyden-Fletcher-Shanno (BFS) matrix update (developed by Broyden [1], Fletcher [3], Goldfarb [5], Greenstadt [6], and Shanno [9]) is used to estimate $H^{(*)}$ in each of the algorithms. Dropping the superscripts, and denoting the current point by x and the subsequent point by x^* , with the corresponding notation g and g^* , and H and H^* , the BFS formula is

$$H^* = (I - \sigma y' / \sigma' y) H (I - \sigma y' / \sigma' y) + \sigma \sigma' / \sigma' y,$$

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where $\sigma = x^* - x$, $y = g^* - g$. This form of the update formula is that suggested by Fletcher [3].

Method 1. This method incorporates an inexact linear search for $\alpha^{(k)}$ at each step. Features of the algorithm include the following.

(i) The initial approximate inverse Hessian matrix proposed by Powell [7] is used, that is, $H^{(0)} = 100 \|\Delta x\| / \|g^{(0)}\| I$, where Δx is the vector of step bounds placed on all variables x_i , $\|\cdot\|$ is the usual L_2 -norm, and I is the $n \times n$ identity matrix.

(ii) In Shanno et al. [11], it has been shown that if α is chosen to satisfy $\sigma'g^* > \sigma'g$, then the positive definiteness of H implies that H^* is also positive definite. Thus α is chosen by the following algorithm. The quadratic extrapolation technique described in [10] is used to bracket the minimum of $f(x)$ along the search direction $s = -Hg$. Davidon's [2] cubic interpolation technique is then applied until a point is found with $f^* < f$ and $\sigma'g^* > \sigma'g$. The point x^* is then accepted as the next point in the sequence.

(iii) Suppose that the step length α along the search direction s obtained by the previous step is "good"; then the distance α^*s^* to be moved from x^* will presumably be about the same as αs . Thus, before applying the above linear-search technique to estimate α , the length of the search vector s is scaled by $\|s^{(k)}\| = \|\sigma^{(k-1)}\|$, $k = 1, \dots, n$. Numerical results indicate that this scaling criterion improves our algorithm for most problems tested.

This program is controlled by setting $\text{MODE} = 1$ in the parameter list of the subroutine MINI. For convenience, we shall denote this algorithm by MINI01. According to the results presented in [12], MINI01 appears preferable to MINI02 on functions where singularity of the true Hessian occurs.

Method 2. In this method we attempt to combine the merits of linear-search methods and step-length methods. The main feature of this algorithm is that linear-search techniques are applied infrequently. Using the notation previously developed, the algorithm can be described by the following seven major steps.

(1) *Create the search direction.* Evaluate $s^* = -H^*g^*$ and scale s^* such that $\|s^*\|$ is equal to $\|\sigma\|$ from the previous iteration for the first n iterations.

(2) *Perform the first extrapolation.* Estimate the first move by calculating

$$\alpha = \max\{-2(f - \bar{f})/s'g, 1\},$$

where \bar{f} is the estimated least value of $f(x)$ provided by the user.

(3) *Evaluate the function value and its gradient.* Set $x^* = x + \alpha s$ and evaluate f^* and g^* . Set $\delta f = f^* - (f + \mu s'g)$, where $\mu = 0.0001$.

(4) *Test for termination of the linear search.* If $\delta f < 0$ and $s'g^* > s'g$, then the search for α is complete. Go to step (7).

(5) *Perform Shanno and Kettler's extrapolation technique.* If $\delta f < 0$ and $s'g^* < s'g$, there exists $\hat{\alpha}$ such that $\hat{\alpha} > \alpha$ and $f(x + \hat{\alpha}s) < f^*$. In this case, apply Shanno and Kettler's [10] extrapolation technique to estimate a new step length, β , that is,

$$\beta = \alpha^2 s'g / (2(f - f^* + \alpha s'g)).$$

Set $x = x^*$ and $\alpha = \beta$, and go to (3). Else go to step (6).

(6) *Perform Davidon's cubic interpolation technique.* If $\delta f \geq 0$ or $s'g^* \geq 0$, the minimum of $f(x)$ along s has been straddled. In this case, apply Davidon's [2] cubic interpolation technique to estimate a new step length, γ , that is,

$$\gamma = (-b + \sqrt{b^2 - 3ac}) / (3a),$$

where

$$\begin{aligned} a &= (2(f - f^*) + \alpha(s'g + s'g^*)) / \alpha^3, \\ b &= (-3(f - f^*) - \alpha(2s'g + s'g^*)) / \alpha^2, \\ c &= s'g. \end{aligned}$$

If $b^2 - 3ac \leq 0$ or $|a| < 10^{-5} |b|$, then apply a quadratic interpolation scheme to estimate γ , that is,

$$\gamma = \alpha^2 s'g / (2(f - f^* + \alpha s'g)).$$

Now set $\alpha = \gamma$ and go to step (3).

(7) *Test for convergence.* If $\|g^*\| \leq \epsilon$, then convergence is assumed. Otherwise update H^* by the BFS formula and repeat the whole process from step (1). Note that this convergence criterion has been found wanting in practice for badly scaled problems. Because it is the criterion used by VMM01 and MINFA, it was necessary to incorporate it in this program in order to obtain reasonable comparisons. A safer test can be incorporated by replacing the statement before statement 400 with

```
IF(DABS(G(I)).GT.EPS*(DABS(X(I)) + .001))CONV = .FALSE.
```

This changes the criterion from an absolute criterion on the gradient to a relative criterion, or the number of positions of accuracy required in the estimates to the minimum.

Remarks

(a) We note that the change in f on an iteration according to Taylor's theorem is approximately $s'g$ when s is small, but much less than $s'g$ in absolute value when the position of the minimum along a line is overestimated. The change in f relative to $s'g$ cannot become arbitrarily small if $(f^* - f)/s'g \geq \mu$, where $0 < \mu \ll 1$ is a preassigned small quantity. For more detail about the choice of this quantity, see Fletcher [3].

(b) On each move of the linear search, α is bounded in such a way that $\alpha s_i \leq \Delta x_i$, $i = 1, 2, \dots, n$, where Δx is the vector of step bounds placed on all variables x_i . The step bounds Δx are introduced for several reasons. First, the initial estimated step length may be far too long, so that much is wasted in finding the desired α . Second, Δx helps in maintaining each component of x within the domain of definition of $f(x)$ so that environmental difficulties such as overflow and underflow problems may be prevented.

(c) In order to expedite the process of linear search, α is only estimated approximately. In fact, the search for α is complete as soon as the conditions given in step (4) are fulfilled. As a consequence, the new step length obtained by applying the extrapolation technique as described in step (5) is restricted to be $\alpha = \min\{\beta, 10\alpha\}$, so that the use of cubic interpolation technique may be avoided.

(d) In real situations, it may happen that the interpolation process gets stuck on one side of the minimum. To overcome this difficulty, the following modification to step (6) is suggested:

$$\alpha = \begin{cases} 0.1 \alpha & \text{if } \gamma \leq 10^{-3}\alpha \\ \gamma & \text{if } 10^{-3}\alpha < \gamma \leq (1 - 10^{-3})\alpha \\ 0.8 \alpha & \text{otherwise.} \end{cases}$$

(e) The use of this algorithm is signified by setting $\text{MODE} = 2$ in the parameter list of the subroutine MINI . For convenience, we shall denote this algorithm by MINI02 . According to the results presented in [12], MINI02 is generally preferable to MINI01 , except when the Hessian has singularities, as noted previously. For the general user, we would recommend use of MINI02 , with a switch to MINI01 if convergence appears "too slow."

Program

The program consists of two modified quasi-Newton minimization techniques (MINI01 , MINI02) which are controlled respectively by setting $\text{MODE} = 1$ or $\text{MODE} = 2$ in the parameter list of the subroutine MINI . This program requires one user-supplied subroutine (CALCFG), which calculates the values of the objective function and its first partial derivatives. The use of the subroutine and the

Table I. Comparison of Numerical Results Obtained by Four Algorithms
 (n = number of function and gradient cells; γ = rank of each algorithm)

	MINIO2		MINIO1		MINFA		VMM01	
	n	γ	n	γ	n	γ	n	γ
BOX								
(5,0)	24	2	40	4	20	1	24	2
(0,0)	19	1	28	4	27	3	21	2
(0,20)	12	1	18	2	21	3	21	3
(2.5,10)	6	1	12	2	19	4	16	3
(5,20)	15	1	21	2	36	4	28	3
Total rank		6		14		16		14
ROSENBROCK								
(1,-1.2)	26	1	42	4	37	3	28	2
(2,-2)	18	1	34	3	39	4	20	2
(-3.635,5.621)	55	2	80	4	39	1	70	3
(6.39,-.221)	64	1	66	2	69	3	74	4
(1.489,-2.547)	24	1	37	3	43	4	28	2
Total rank		6		16		15		13
BEALE								
(1,-1.2)	13	1	23	4	15	2	16	3
(2,-2)	18	1	26	4	25	3	24	2
(0,0)	13	1	22	3	22	3	13	1
(5,3)	36	4	35	3	25	2	11	1
(10,10)	74	2	84	4	77	3	18	1
Total rank		9		18		14		9
WEIBULL								
(5,.15,2.5)	44	1	84	2	144	3	(3)	4+F
(250,.3,5)	63	1	94	4	81	2	93	3
(100,3,12.5)	48	1	66	2	74	3	(4)	4+F
Total rank		3		8		8		11+2F
WOOD								
(-3,-1,-3,-1)	99	2	91	1	113	3	132	4
(-3,1,-3,1)	98	2	78	1	100	3	153	4
(-1.2,1,-1.2,1)	84	2	54	1	112	4	95	3
(-1.2,1,1.2,1)	45	1	95	4	70	3	64	2
Total rank		7		7		13		13
POWELL								
(-3,-1,0,1)	39	1	57	3	(1)	4+F	47	2
CREGG								
(1,2,2,2)	45	1	79	2	(2)	3+F	(5)	31+F
Overall Total Rank		33		68		71+2F		61+3F

- (1) = underflow occurred at the 73th iteration.
 (2) = underflow occurred at the 123th iteration.
 (3) = the program exited with IEXIT=4, see Fletcher (4).
 (4) = underflow occurred at the 2nd iteration.
 (5) = the program exited with FEXIT=4, see Fletcher (4).

meaning of the parameters are described in the comments at the beginning of subroutine MINI. All communication between the main program and subroutines is achieved through the subroutine argument lists. An iteration is defined as the calculations required to select a new point which yields a lower function value than that of the previous one.

Test Results

These two programs (MINIO1, MINIO2) have been tested over a wide range of test functions in [12]. Compared with the subroutine (VMM01) given by Fletcher [3] and the subroutine (MINFA) as suggested by Powell [7], results show that MINIO2 is the best algorithm, while MINIO1 places second. For user comparisons,

we list some computational results of these two programs for the following four standard test functions.

(1) Box's function, defined by

$$f(x_1, x_2) = \sum_{i=1}^{10} \left| (e^{-x_1 t_i} - e^{-x_2 t_i}) - (e^{-t_i} - e^{-10 t_i}) \right|^2,$$

where t_i ranges from .1 to 1 in steps of .1. For MINI01 and MINI02, the step bounds are set to be $\text{STP}(i) = 3$, $i = 1, 2$, and $\Delta^{(0)} = 3$ for MINFA.

(2) Rosenbrock's function, defined by

$$f(x_1, x_2) = 100(x_2 - x_1^2)^2 + (1 - x_1)^2,$$

with $\text{STP}(i) = 3$, $i = 1, 2$, and $\Delta^{(0)} = 3$ for MINFA.

(3) Wood's function, defined by

$$\begin{aligned} f(x_1, x_2, x_3, x_4) = & 100(x_2 - x_1^2)^2 + (1 - x_1)^2 + 90(x_4 - x_3^2)^2 \\ & + (1 - x_3)^2 + 10.1[(x_2 - 1)^2 + (x_4 - 1)^2] \\ & + 19.8(x_2 - 1)(x_4 - 1), \end{aligned}$$

with $\text{STP}(i) = 10$, $i = 1, 2, 3, 4$, and $\Delta^{(0)} = 10$ for MINFA.

(4) Weibull's function, defined by

$$f(x_1, x_2, x_3) = \sum_{i=1}^{99} [\exp((x_3 - t_i)x^2/x_1) - y_i]^2,$$

where the y_i and t_i are perfect data generated for the 99 points corresponding to $y = .01$ to $.99$ in steps of $.01$, for the values $x_1 = 50$, $x_2 = 1.5$, $x_3 = 25$, and $\text{STP}(1) = 300$, $\text{STP}(2) = 3$, $\text{STP}(3) = 30$, and $\Delta^{(0)} = 3$ for MINFA.

(5) Beale's function, defined by

$$f(x_1, x_2) = [1.5 - x_1(1 - x_2)]^2 + [2.25 - x_1(1 - x_2^2)]^2 + [2.625 - x_1(1 - x_2^3)]^2,$$

with $\Delta^{(0)} = 3$ for MINFA and $|\Delta x_i| \leq 3$, $i = 1, 2$, for MINI01 and MINI02.

(6) Powell's function, defined by

$$f(x_1, x_2, x_3, x_4) = (x_1 + 10x_2)^2 + 5(x_3 - x_4)^2 + (x_2 - 2x_3)^4 + 10(x_1 - x_4)^4,$$

with $\Delta^{(0)} = 3$ for MINFA, and $|\Delta x_i| \leq 3$, $i = 1, 2, 3, 4$, for MINI01 and MINI02.

(7) Cregg's function, defined by

$$f(x_1, x_2, x_3, x_4) = (e^{x_1} - x_2)^4 + 100(x_2 - x_3)^6 + \tan^4(x_3 - x_4) + x_1^8 + (x_4 - 1)^2,$$

with $\Delta^{(0)} = 3$ for MINFA, and $|\Delta x_i| \leq 3$, $i = 1, 2, 3, 4$, for MINI01 and MINI02.

All tests were run in double precision on an IBM 370/165. In all cases, convergence was assumed to have occurred when $|\partial f/\partial x_i| \leq 10^{-5}$ for each i . Each test function was started from various initial estimates and their results are listed in Table I. The table does not differentiate function from gradient calls, because all programs tested calculate both each time the function subroutine is called. For a more complete evaluation, encompassing overhead and execution time, see [12]. We note here only that more careful evaluation leaves relative rankings unchanged.

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ALGORITHM

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SUBROUTINE MINI(N, X, F, G, H, M, IH, STP, EPS, FEST, MAXF, MIN 10
* MODE, ITER, IFUN, IERR, CALCFG) MIN 20
C PURPOSE. THIS SUBROUTINE FINDS A MINIMUM OF THE FUNCTION F(X) MIN 30
C ----- WITHIN THE ACCURACY NORM(G(X)) .LE. EPS, WHERE G(X) IS MIN 40
C THE GRADIENT OF F(X) AND EPS IS THE REQUIRED ACCURACY MIN 50
C PROVIDED BY THE USER. MIN 60
C THIS SUBROUTINE ADOPTS THE QUASI-NEWTON METHOD AS MIN 70
C SUGGESTED BY SHANNO, D.F. AND PHUA, K.H. (1974). MIN 80
C DESCRIPTION OF PARAMETERS IN THE ARGUMENT LIST MIN 90
C ----- MIN 100
C N NUMBER OF VARIABLES MIN 110
C X STORES THE CURRENT APPROXIMATION TO THE VARIABLES X. AN MIN 120
C INITIAL APPROXIMATION SHOULD BE PROVIDED ON ENTRY, WHICH MIN 130
C WILL BE REPLACED BY THE BEST ESTIMATE OBTAINED ON EXIT. MIN 140
C F ON RETURN, F CONTAINS THE CORRESPONDING VALUE OF F(X). MIN 150
C G ON RETURN, G IS THE CORRESPONDING GRADIENT VECTOR G(X). MIN 160
C H USED TO STORE THE APPROXIMATED INVERSE HESSIAN MATRIX AS MIN 170
C ROWS OF ITS UPPER TRIANGLE. MIN 180
C M = N*(N+1) / 2, THE LENGTH OF THE VECTOR H. MIN 190
C IH INDICATOR CHOOSING THE INITIAL APPROXIMATION TO THE INVERSE MIN 200
C HESSIAN MATRIX, MIN 210
C = 0 THE MINI SUBROUTINE SETS H INITIALLY. MIN 220
C = 1 IT IS ASSUMED THAT A GOOD ESTIMATE OF THIS MATRIX IS MIN 230
C PROVIDED, FOR INSTANCE FROM PREVIOUS CALCULATIONS. MIN 240
C STP THE VECTOR CONTAINS THE MAXIMUM ALLOWABLE STEP-LENGTH ON MIN 250
C EACH COMPONENT OF X. THE USER MUST PROVIDE THIS VECTOR. MIN 260
C PROPER CHOICE OF THE COMPONENTS STP(I) IS IMPORTANT FOR MIN 270
C EFFICIENT EXECUTION. IN GENERAL, CHOOSING STP(I) TOO LARGE MIN 280
C IS PREFERABLE TO CHOOSING STP(I) TOO SMALL. THUS STP(I) MIN 290
C SHOULD IDEALLY BE CHOSEN TO LIMIT SEARCH TO THE SMALLEST MIN 300
C REGION ABOUT X(1) IN WHICH THE USER HAS CONFIDENCE THE TRUE MIN 310
C OPTIMUM LIES. IF THE USER HAS NO FEEL FOR THE LOCATION OF MIN 320
C THE TRUE OPTIMUM, STP(I) SHOULD BE CHOSEN TO BE LARGE, MIN 330
C REFLECTING THE USERS VIEW THAT THE OPTIMUM MAY BE FAR FROM MIN 340
C THE INITIAL ESTIMATES. MIN 350
C EPS ACCURACY REQUIRED IN EACH ELEMENT OF THE GRADIENT VECTOR MIN 360
C G(X). IF THE USER HAS NO FEEL FOR THE DESIRED ACCURACY OF MIN 370
C HIS SOLUTION, EPS = 10**(-5) IS SUGGESTED. MIN 380
C FEST A LOWER BOUND ON THE VALUE OF F PROVIDED BY THE USER. MIN 390
C MAXF MAXIMUM ALLOWABLE FUNCTION EVALUATIONS, PROVIDED BY THE MIN 400
C USER. MAXF ACTS SOLELY AS A DEVICE TO LIMIT EXECUTION TIME MIN 410
C IN THE EVENT OF SLOW CONVERGENCE. IF THE USER HAS NO FEEL MIN 420
C FOR EXECUTION TIME AS A FUNCTION OF TIME/ITERATION, MIN 430
C MAXF = 200 IS SUGGESTED. MIN 440
C MODE ROUTINE SELECTOR PROVIDED BY THE USER, MIN 450
C = 1 MEANS THE ROUTINE MINI01 WILL BE USED. MIN 460
C = 2 MEANS THE ROUTINE MINI02 WILL BE USED. MIN 470
C MODE = 2 IS SUGGESTED UNLESS CONVERGENCE IS UNACCEPTABLY MIN 480
C SLOW. IN THIS CASE, MODE = 1 SHOULD BE TRIED. MIN 490
C ITER NUMBER OF ITERATIONS USED. MIN 500
C IFUN NUMBER OF FUNCTION EVALUATIONS USED. MIN 510
C IERR ERROR MESSAGE RETURNED, MIN 520
C = 0 MEANS NO ERROR ENCOUNTERED. MIN 530
C = 1 MEANS IT REQUIRES MORE FUNCTION EVALUATIONS. MIN 540
C = 2 MEANS INCREMENT IN X IS TOO SMALL, PROGRAM STUCK, MIN 550
C POSSIBLY CAUSED BY 1) H SINGULAR, 2) EPS SET TOO SMALL MIN 560
C (SEEK EXPERT FOR ADVICE). MIN 570
C = 3 MEANS BOTH THE SEARCH DIRECTION AND ITS REVERSE MIN 580
C DIRECTION ARE NOT DOWN-HILL, POSSIBLY CAUSED BY G MIN 590
C PROGRAMMED INCORRECTLY. MIN 600
C = 4 MEANS THAT FUNCTION VALUE LESS THAN FEST DETECTED. MIN 610
C REMARKS MIN 620
C ----- MIN 630
C 1) THE USER MUST PROVIDE A SUBROUTINE MIN 640
C SUBROUTINE CALCFG (N, X, F, G) MIN 650
C DOUBLE PRECISION X(N), G(N), F MIN 660
C WHICH, GIVEN A VECTOR X IN X(1), X(2), ..., X(N), CALCULATES THE MIN 670
C FUNCTION F(X) AND THE GRADIENT G(X) AND PLACES THEM IN F AND MIN 680
C G(1), G(2), ..., G(N) RESPECTIVELY. MIN 690
C 2) THE VECTOR D(10) IS USED TO HOLD THE CURRENT SEARCH DIRECTION, MIN 700
C WHILE XX(10) AND GG(10) ARE USED TO STORE THE PREVIOUS VALUES MIN 710
C OF X AND G. THUS IF THE VALUE OF N IS GREATER THAN 10, THEN MIN 720
C THEIR DIMENSIONS MUST BE CHANGED ACCORDINGLY. MIN 730

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      DOUBLE PRECISION F, FEST, EPS, X(N), G(N), STP(N), H(M)
      DOUBLE PRECISION D(10), XX(10), GG(10), EPS2, DUM, STEP, SUM,
      * G1, G2
      DOUBLE PRECISION STPMAX, ALPHA, AL, BL, FG, FL, FM, GM, AA,
      * BB, CC
      DOUBLE PRECISION SS, DY, YHY, C1, C2, DSQRT, DABS
      LOGICAL CONV
C DO INITIALIZATION OF PARAMETERS, STEP-BOUND AND UPDATE
      CALL CALCFG(N, X, F, G)
      IFUN = 1
      ITER = 0
      IERR = 0
      EPS2 = EPS*EPS
C CALCULATE THE INITIAL STEP-BOUND
      DUM = 0.
      STEP = 0.
      DO 10 I=1,N
          DUM = DUM + G(I)*G(I)
          STEP = STEP + STP(I)*STP(I)
      10 CONTINUE
      IF (DUM.EQ.0.) RETURN
C SET UP THE INITIAL UPDATE
      IF (IH.EQ.0) GO TO 20
      GO TO 50
C THEN
      20 DUM = 100.*DSQRT(STEP/DUM)
      IJ = 1
      DO 40 I=1,N
          DO 30 J=I,N
              H(IJ) = 0.
              IF (I.EQ.J) H(IJ) = DUM
              IJ = IJ + 1
          30 CONTINUE
      40 CONTINUE
      50 CONTINUE
C REPEAT THE FOLLOWING 5 STEPS UNTIL EXIT FROM 4TH STEP
C *****
C * STEP 1. SET UP THE SEARCH DIRECTION AND SCALE ITS LENGTH *
C *****
      60 SUM = 0.
      DO 110 I=1,N
          DUM = 0.
          IJ = I
          IF (I.GT.1) GO TO 70
          GO TO 90
C THEN
      70 II = I - 1
      DO 80 J=1,II
          DUM = DUM - H(IJ)*G(J)
          IJ = IJ + N - J
      80 CONTINUE
      90 CONTINUE
      DO 100 J=I,N
          DUM = DUM - H(IJ)*G(J)
          IJ = IJ + 1
      100 CONTINUE
          D(I) = DUM
          XX(I) = X(I)
          GG(I) = G(I)
          SUM = SUM + DUM*DUM
      110 CONTINUE
          IF ((ITER.GE.1) .AND. (ITER.LE.N)) GO TO 120
          GO TO 140
C THEN - SCALE THE LENGTH OF THE SEARCH VECTOR
      120 DUM = DSQRT(STEP/SUM)
      DO 130 I=1,N
          D(I) = D(I)*DUM
      130 CONTINUE
      140 CONTINUE
C *****
C * STEP 2. TEST IF THE SEARCH DIRECTION IS DOWNHILL *
C *****
      DO 170 J=1,2
          G1 = 0.
          DO 150 I=1,N
              G1 = G1 + D(I)*G(I)
          150 CONTINUE
          IF (G1.LT.0.) GO TO 180
          DO 160 I=1,N
              D(I) = -D(I)
          160 CONTINUE
      170 CONTINUE
      180 IF (G1.GT.0.) GO TO 190
          GO TO 200
C THEN
      190 IERR = 3
C .....EKKOR EXIT - BOTH D AND -D ARE NOT DOWNHILL
      RETURN
      200 CONTINUE
C FIND OUT THE MAXIMUM ALLOWABLE STEP LENGTH, STPMAX
      STPMAX = 1.0D+30

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```

DO 210 I=1,N
  IF (DABS(D(I)).LT.1.D-30) GO TO 210
  DUM = DABS(STP(I)/D(I))
  IF (STPMAX.GT.DUM) STPMAX = .9*DUM
210 CONTINUE
C *****
C * STEP 3. PERFORM THE LINEAR SEARCH *
C *****
  ALPHA = 0.
  FG = G1
  AL = 1.
  IF (MODE.EQ.2) AL = 2.*(FEST-F)/G1
  IF (AL.GT.1.) AL = 1.
  IF (AL.GT.STPMAX) AL = 0.8*STPMAX
  IF (AL.LT.0.) GO TO 220
  GO TO 230
C THEN
220 IERR = 4
C .....ERROR EXIT - F .LT. FEST PROVIDED BY THE USER
  RETURN
230 CONTINUE
C REPEAT THE FOLLOWING EXTRAPOLATION PROCESS
240 F1 = F
  ALPHA = ALPHA + AL
  DO 250 I=1,N
    X(I) = XX(I) + ALPHA*D(I)
250 CONTINUE
  CALL CALCFG(N, X, F, G)
  IFUN = IFUN + 1
  IF (IFUN.GT.MAXF) GO TO 260
  GO TO 270
C THEN
260 IERR = 1
C .....ERROR EXIT - EXCEED THE MAXIMUM ALLOWED FUNCTION CALLS
  RETURN
270 CONTINUE
  G2 = 0.0
  DO 280 I=1,N
    G2 = G2 + G(I)*D(I)
280 CONTINUE
  CONV = (G2.GT.FG) .AND. (F1.GT.(F+.0001*G1))
C ...EXIT THE EXTRAPOLATION LOOP, IF THE MINIMUM HAS BEEN FOUND
  IF (CONV .AND. (MODE.EQ.2)) GO TO 290
C ...EXIT THE EXTRAPOLATION LOOP, IF THE MINIMUM IS STRADDLED
  IF ((F.GT.F1) .OR. (G2.GT.0.)) GO TO 290
  BL = AL
  AL = 0.5*G1*BL*BL/(F1-F+BL*G1)
  IF ((MODE.EQ.2) .AND. (AL.LT.BL)) AL = 2.*(FEST-F)/G1
  IF (AL.GT.(10.*BL)) AL = 10.*BL
  IF (AL.GT.STPMAX) AL = 0.8*STPMAX
  IF (AL.LT.(1.001*BL)) AL = BL + BL
  G1 = G2
C *GO BACK TO THE EXTRAPOLATION PROCESS
  GO TO 240
290 CONTINUE
  IF (.NOT.(CONV) .OR. (MODE.EQ.1)) GO TO 300
  GO TO 390
C THEN
C REPEAT THE FOLLOWING INTERPOLATION PROCESS
300 BL = AL
  AA = (G1+G2+2.*(F1-F)/AL)/(AL*AL)
  BB = (G2-3.*AA*AL*AL-G1)/(AL+AL)
  CC = BB*BB - 3.*AA*G1
  DUM = DABS(AA)*1.0D+05 - DABS(BB)
  AL = (.5*G1*BL*BL)/(F1-F+BL*G1)
  IF ((CC.GT.0.) .AND. (DUM.GT.0.)) AL = (-BB+DSQRT(CC))/(3.*AA)
  FM = F
  GM = G2
  IF (AL.LE.(.001*BL)) AL = .1*BL
  IF (AL.GT.(.999*BL)) AL = .8*BL
  ALPHA = ALPHA - BL + AL
  DO 310 I=1,N
    X(I) = XX(I) + ALPHA*D(I)
310 CONTINUE
  CALL CALCFG(N, X, F, G)
  IFUN = IFUN + 1
  IF (IFUN.GT.MAXF) GO TO 320
  GO TO 330
C THEN
320 IERR = 1
C .....ERROR EXIT - EXCEED THE MAXIMUM ALLOWED FUNCTION
C CALLS
  RETURN
330 CONTINUE
  G2 = 0.
  DO 340 I=1,N
    G2 = G2 + D(I)*G(I)
340 CONTINUE
  CONV = (G2.GT.FG) .AND. (F1.GT.(F+.0001*G1))
C ...EXIT THE INTERPOLATION LOOP, IF THE MINIMUM IS FOUND
  IF (CONV) GO TO 380

```

```

MIN 1650
MIN 1660
MIN 1670
MIN 1680
MIN 1690
MIN 1700
MIN 1710
MIN 1720
MIN 1730
MIN 1740
MIN 1750
MIN 1760
MIN 1770
MIN 1780
MIN 1790
MIN 1800
MIN 1810
MIN 1820
MIN 1830
MIN 1840
MIN 1850
MIN 1860
MIN 1870
MIN 1880
MIN 1890
MIN 1900
MIN 1910
MIN 1920
MIN 1930
MIN 1940
MIN 1950
MIN 1960
MIN 1970
MIN 1980
MIN 1990
MIN 2000
MIN 2010
MIN 2020
MIN 2030
MIN 2040
MIN 2050
MIN 2060
MIN 2070
MIN 2080
MIN 2090
MIN 2100
MIN 2110
MIN 2120
MIN 2130
MIN 2140
MIN 2150
MIN 2160
MIN 2170
MIN 2180
MIN 2190
MIN 2200
MIN 2210
MIN 2220
MIN 2230
MIN 2240
MIN 2250
MIN 2260
MIN 2270
MIN 2280
MIN 2290
MIN 2300
MIN 2310
MIN 2320
MIN 2330
MIN 2340
MIN 2350
MIN 2360
MIN 2370
MIN 2380
MIN 2390
MIN 2400
MIN 2410
MIN 2420
MIN 2430
MIN 2440
MIN 2450
MIN 2460
MIN 2470
MIN 2480
MIN 2490
MIN 2500
MIN 2510
MIN 2520
MIN 2530
MIN 2540
MIN 2550

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C ...EXIT THE INTERPOLATION LOOP, IF STEP-SIZE IS TOO SMALL
  IF ((G2.GT.FG) .AND. (BL.LE.EPS2)) GO TO 380
  SS = 0.
  DO 350 I=1,N
    SS = SS + DABS(G(I)*D(I))
  350 CONTINUE
  SS = SS/(100000.*FLOAT(N))
C ...EXIT THE INTERPOLATION LOOP, IF TWO POINTS ARE CLOSED
  IF ((G2.LT.SS) .AND. (G2.GT.FG)) GO TO 380
  IF ((G2.LT.0.) .AND. (F1.GT.F)) GO TO 360
  GO TO 370
C THEN - ADVANCE TO THE LOWER POINT
  360 G1 = G2
  G2 = GM
  F1 = F
  F = FM
  AL = BL - AL
  ALPHA = ALPHA + AL
  370 CONTINUE
C *GO BACK TO THE INTERPOLATION PROCESS
  GO TO 300
  380 CONTINUE
  390 CONTINUE
C *****
C * STEP 4. TEST FOR CONVERGENCE *
C *****
  CONV = .TRUE.
  STEP = 0.
  DO 400 I=1,N
    D(I) = X(I) - XX(I)
    STEP = STEP + D(I)*D(I)
    IF ((G(I)*G(I)).GT.EPS2) CONV = .FALSE.
  400 CONTINUE
C .....NORMAL EXIT
  IF (CONV) RETURN
  IF (STEP.LE.(EPS2*EPS2)) GO TO 410
  GO TO 420
C THEN
  410 IERR = 2
C .....ERROR EXIT - THE STEP SIZE IS TOO SMALL
  RETURN
  420 CONTINUE
C *****
C * STEP 5. UPDATE THE APPROXIMATED INVERSE HESSIAN MATRIX *
C *****
  DY = 0.
  YHY = 0.
  DO 430 I=1,N
    GG(I) = G(I) - GG(I)
    DY = DY + D(I)*GG(I)
  430 CONTINUE
  DO 480 I=1,N
    DUM = 0.
    IJ = I
    IF (I.GT.1) GO TO 440
    GO TO 460
  C THEN
  440 II = I - 1
  DO 450 J=1,II
    DUM = DUM + H(IJ)*GG(J)
    IJ = IJ + N - J
  450 CONTINUE
  460 CONTINUE
  DO 470 J=I,N
    DUM = DUM + H(IJ)*GG(J)
    IJ = IJ + 1
  470 CONTINUE
  YHY = YHY + DUM*GG(I)
  XX(I) = DUM
  480 CONTINUE
  C1 = 1. + YHY/DY
  DO 490 I=1,N
    GG(I) = C1*D(I) - XX(I)
  490 CONTINUE
  IJ = 1
  DO 510 I=1,N
    C1 = D(I)/DY
    C2 = XX(I)/DY
    DO 500 J=I,N
      H(IJ) = H(IJ) + C1*GG(J) - C2*D(J)
      IJ = IJ + 1
  500 CONTINUE
  510 CONTINUE
  ITER = ITER + 1
C *STEP 5 NOW COMPLETE - GO BACK TO STEP 1
  GO TO 60
C *END OF REPEAT BLOCK
C CONTINUE
  END

```

MIN 2560
MIN 2570
MIN 2580
MIN 2590
MIN 2600
MIN 2610
MIN 2620
MIN 2630
MIN 2640
MIN 2650
MIN 2660
MIN 2670
MIN 2680
MIN 2690
MIN 2700
MIN 2710
MIN 2720
MIN 2730
MIN 2740
MIN 2750
MIN 2760
MIN 2770
MIN 2780
MIN 2790
MIN 2800
MIN 2810
MIN 2820
MIN 2830
MIN 2840
MIN 2850
MIN 2860
MIN 2870
MIN 2880
MIN 2890
MIN 2900
MIN 2910
MIN 2920
MIN 2930
MIN 2940
MIN 2950
MIN 2960
MIN 2970
MIN 2980
MIN 2990
MIN 3000
MIN 3010
MIN 3020
MIN 3030
MIN 3040
MIN 3050
MIN 3060
MIN 3070
MIN 3080
MIN 3090
MIN 3100
MIN 3110
MIN 3120
MIN 3130
MIN 3140
MIN 3150
MIN 3160
MIN 3170
MIN 3180
MIN 3190
MIN 3200
MIN 3210
MIN 3220
MIN 3230
MIN 3240
MIN 3250
MIN 3260
MIN 3270
MIN 3280
MIN 3290
MIN 3300
MIN 3310
MIN 3320
MIN 3330
MIN 3340
MIN 3350
MIN 3360
MIN 3370
MIN 3380
MIN 3390
MIN 3400
MIN 3410
MIN 3420
MIN 3430
MIN 3440

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REMARK ON ALGORITHM 500

Minimization of Unconstrained Multivariate Functions [E4]

[D.F. Shanno and K.H. Phua, *ACM Trans. Math. Software* 2, 1 (March 1976), 87–94]

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The arrays STP and H used as arguments to this algorithm must be declared double precision, as declared in the body (but not the published heading) of the subroutine MINI.

MINI was used to determine a best $l_{1.5}$ approximation by the approximating function

$$F(A, x) = a_1 + a_2x + a_3 \exp(a_4x)$$

on $\{0, 1/10, \dots, 9/10, 1\}$ to the function $f(x) = x + \exp(x) + \exp(2x)/100$. This is a minimization problem with four parameters, a_1, a_2, a_3, a_4 . When MINI was run initially with arguments STP single precision of dimension 4 and H single precision of dimension 10 on a CDC Cyber 73, the program hung with an error mode 2 (infinite operand). With the above change (to double precision) the program produced the approximation obtained independently by several other methods.

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Minimization of Unconstrained Multivariate Functions [E4] [D.F. Shanno and K.H. Phua, *ACM Trans. Math. Softw.* 2, 1 (March 1976), 87–94.]

D.F. Shanno and K.H. Phua [Received 3 January 1979; revised 5 June 1979; accepted 26 December 1979]

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DESCRIPTION

1. Purpose

This subroutine finds a local minimizer of a nonlinear function of n variables $f(x)$ where $x = (x_1, \dots, x_n)$, $n \geq 1$ can be any real numbers. This algorithm is meant to supersede the algorithm documented in [11].

2. Methods

The subroutine incorporates two nonlinear optimization methods, a conjugate gradient algorithm and a variable metric algorithm, with the choice of method left to the user.

The conjugate gradient algorithm is the Beale restarted memoryless variable metric algorithm documented in Shanno [7]. This method requires approximately $7n$ double-precision words of working storage to be provided by the user. The variable metric method is the BFGS algorithm with initial scaling documented in Shanno and Phua [10], and required approximately $n^2/2 + 11n/2$ double-precision words of working storage.

Whichever method is chosen, the same linear search technique is used for both methods, with two differences. The basic linear search uses Davidon's cubic interpolation to find a step length α , which satisfies

$$f(x + \alpha d) < f(x) + \alpha d'g(x) 0.0001, \quad (1)$$

where d is the chosen search direction, $g(x) = \nabla f(x)$, the gradient of f at x , and $d'g(x)$, the directional derivative of $f(x)$ at x along d , is always a negative number. In addition, α must satisfy

$$|d'g(x + \alpha d)/d'g(x)| < 0.9. \quad (2)$$

The convergence of the BFGS variable metric algorithm under the conditions (1) and (2) has been studied by Powell [6], while the convergence of the conjugate gradient method has been studied by Shanno [8].

The major difference between the two methods insofar as the linear search is concerned is that if the first trial α satisfies (1) and (2), it is accepted if a variable metric method is used, but at least two trial α 's are required before accepting an α satisfying (1) and (2) for the conjugate gradient method. Reasons for requiring this are detailed in Shanno [9].

The second difference between the two methods is that for a variable metric method, $\alpha = 1$ is always the initial α tried, while for the conjugate gradient method, $\alpha = 1$ is tried only for restart iterations, whereas for nonrestart iterations the initial α at step $k + 1$, denoted by α_{k+1} , is chosen to be $\alpha_{k+1} = \alpha_k d'_k g_k / d'_{k+1} g_{k+1}$, where d_k and g_k and d_{k+1} and g_{k+1} are the search vectors and gradients at the k th and $k + 1$ st points, respectively.

The linear search contains safeguards to ensure that the search procedure cannot become stuck or attempt to move past a local maximum to a more distant local minimum.

Convergence is determined to have occurred when $\|g\| < \epsilon \max(1, \|x\|)$, where $\|\cdot\|$ is the Euclidean norm and ϵ is user supplied.

3. Test Results

Table I contains test results for a variety of test problems for both the BFGS and conjugate gradient method. The test functions are Wood's function for the listed initial estimates, the extended Rosenbrock function documented in [7], Brodlić's [1] variable-dimensional Watson function, Oren's [5] power function, Powell's four-dimensional function [6], Fletcher and Powell's trigonometric functions with initial estimates as in [2], and for various dimensions, the Mancino function with initial estimates as documented in [7], Moré, Garbow, and Hillstom's [4] boundary-value problem, and Toint's variation on Broyden's function [12]. The Wood and Powell functions are documented in [3].

In the table, ITER represents the number of search directions calculated, while IFUN represents the number of function and gradient evaluations that were performed. In all cases, $\epsilon = 10^{-5}$ was used, except for the boundary value problem, where $\epsilon = 10^{-4}$.

With the exception of the power function, on which the BFGS does not perform well for reasons documented in [10], it is clear from Table I that the variable metric method is quite a bit more efficient in terms of function and gradient calls, primarily due to the fact that the first trial α can be accepted, while at least two trial α 's per iteration must be tried by the conjugate gradient algorithm.

In terms of execution time, however, the issue is not so clear-cut. On a DEC-10 computer, the Broyden-Toint function with $n = 30$ took 6.49 CPU seconds for the BFGS, while the conjugate gradient method took 5.32 due to the overhead of updating the approximate Hessian at each step. However, for the 30-variable Mancino function, the BFGS took 28.12 CPU seconds, while the conjugate gradient method took 72.31. As one would expect, evaluations of the Mancino function are quite expensive, while the Broyden-Toint function evaluations are quite inexpensive.

Thus for large problems where space limitations do not preclude using the BFGS algorithm, users are urged to experiment to determine the most efficient algorithm for a particular problem. For small problems, we recommend the BFGS method, while for very large problems, memory considerations generally mandate using the conjugate gradient algorithm.

Table I

	BFGS		Conjugate Gradient	
	ITER	IFUN	ITER	IFUN
WOOD ($n = 4$)				
-3, -1, -3, -1	36	43	48	106
-3, 1, -3, 1	91	114	90	210
-1.2, 1, -1.2, 1	87	107	77	181
-1.2, 1, 1.2, 1	48	57	46	100
EROSÉN				
-1.2, 1, 1, 1, 1 ($n = 5$)	117	150	132	278
-1, ..., -1 ($n = 10$)	674	893	946	1940
WATSON				
0, ..., 0 ($n = 5$)	37	39	34	69
0, ..., 0 ($n = 10$)	92	95	179	360
POWER				
1, ..., 1 ($n = 20$)	280	281	16	33
1, ..., 1 ($n = 50$)	539	540	30	61
POWELL ($n = 4$)				
-3, -1, 0, 1	48	49	28	57
TRIG				
($n = 5$)	20	22	20	41
($n = 10$)	35	37	44	89
($n = 15$)	57	59	100	201
MANCINO				
($n = 10$)	9	10	12	28
($n = 20$)	10	14	14	33
($n = 30$)	11	17	18	49
BOUNDARY VALUE				
($n = 10$)	28	30	25	51
($n = 20$)	56	58	48	97
($n = 30$)	86	88	121	243
BROYDEN-TOINT				
-1, ..., -1 ($n = 10$)	27	28	23	47
-1, ..., -1 ($n = 20$)	36	37	36	73
-1, ..., -1 ($n = 30$)	47	48	46	93

ACKNOWLEDGMENT

The authors are deeply indebted to Professor M.J.D. Powell for many useful suggestions concerning the linear search method adopted and for illustrating several adopted methods for coding optimization routines.

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ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```

SUBROUTINE CONMIN(N,X,F,G,IFUN,ITER,EPS,NFLAG,MXFUN,W,
  IOUT,MDIM,IDEV,ACC,NMETH)
C
C PURPOSE:  SUBROUTINE CONMIN MINIMIZES AN UNCONSTRAINED NONLINEAR
C           SCALAR VALUED FUNCTION OF A VECTOR VARIABLE X
C           EITHER BY THE BFGS VARIABLE METRIC ALGORITHM OR BY A
C           BEALE RESTARTED CONJUGATE GRADIENT ALGORITHM.
C
C USAGE:    CALL CONMIN(N,X,F,G,IFUN,ITER,EPS,NFLAG,MXFUN,W,
C           IOUT,MDIM,IDEV,ACC,NMETH)
C
C PARAMETERS: N    THE NUMBER OF VARIABLES IN THE FUNCTION TO
C                 BE MINIMIZED.
C                 X    THE VECTOR CONTAINING THE CURRENT ESTIMATE TO
C                 THE MINIMIZER. ON ENTRY TO CONMIN,X MUST CONTAIN
C                 AN INITIAL ESTIMATE SUPPLIED BY THE USER.
C                 ON EXITING,X WILL HOLD THE BEST ESTIMATE TO THE
C                 MINIMIZER OBTAINED BY CONMIN. X MUST BE DOUBLE
C                 PRECISIONED AND DIMENSIONED N.
C                 F    ON EXITING FROM CONMIN,F WILL CONTAIN THE LOWEST
C                 VALUE OF THE OBJECT FUNCTION OBTAINED.
C                 F IS DOUBLE PRECISIONED.
C                 G    ON EXITING FROM CONMIN,G WILL CONTAIN THE
C                 ELEMENTS OF THE GRADIENT OF F EVALUATED AT THE
C                 POINT CONTAINED IN X. G MUST BE DOUBLE
C                 PRECISIONED AND DIMENSIONED N.
C                 IFUN  UPON EXITING FROM CONMIN,IFUN CONTAINS THE
C                 NUMBER OF TIMES THE FUNCTION AND GRADIENT
C                 HAVE BEEN EVALUATED.
C                 ITER  UPON EXITING FROM CONMIN,ITER CONTAINS THE
C                 TOTAL NUMBER OF SEARCH DIRECTIONS CALCULATED
C                 TO OBTAIN THE CURRENT ESTIMATE TO THE MINIZER.
C                 EPS   EPS IS THE USER SUPPLIED CONVERGENCE PARAMETER.
C                 CONVERGENCE OCCURS WHEN THE NORM OF THE GRADIENT
C                 IS LESS THAN OR EQUAL TO EPS TIMES THE MAXIMUM
C                 OF ONE AND THE NORM OF THE VECTOR X. EPS
C                 MUST BE DOUBLE PRECISIONED.
C                 NFLAG UPON EXITING FROM CONMIN,NFLAG STATES WHICH
C                 CONDITION CAUSED THE EXIT.
C                 IF NFLAG=0, THE ALGORITHM HAS CONVERGED.
C                 IF NFLAG=1, THE MAXIMUM NUMBER OF FUNCTION
C                 EVALUATIONS HAVE BEEN USED.
C                 IF NFLAG=2, THE LINEAR SEARCH HAS FAILED TO
C                 IMPROVE THE FUNCTION VALUE. THIS IS THE
C                 USUAL EXIT IF EITHER THE FUNCTION OR THE
C                 GRADIENT IS INCORRECTLY CODED.
C                 IF NFLAG=3, THE SEARCH VECTOR WAS NOT
C                 A DESCENT DIRECTION. THIS CAN ONLY BE CAUSED
C                 BY ROUND OFF,AND MAY SUGGEST THAT THE
C                 CONVERGENCE CRITERION IS TOO STRICT.
C                 MXFUN MXFUN IS THE USER SUPPLIED MAXIMUM NUMBER OF
C                 FUNCTION AND GRADIENT CALLS THAT CONMIN WILL
C                 BE ALLOWED TO MAKE.
C                 W     W IS A VECTOR OF WORKING STORAGE.IF NMETH=0,
C                 W MUST BE DIMENSIONED 5*N+2. IF NMETH=1,
C                 W MUST BE DIMENSIONED N*(N+7)/2. IN BOTH CASES,
C                 W MUST BE DOUBLE PRECISIONED.
C                 IOUT  IOUT IS A USER SUPPLIED OUTPUT PARAMETER.
C                 IF IOUT = 0, THERE IS NO PRINTED OUTPUT FROM
C                 CONMIN. IF IOUT J 0, THE VALUE OF F AND THE
C                 NORM OF THE GRADIENT SQUARED,AS WELL AS ITER
C                 AND IFUN,ARE WRITTEN EVERY IOUT ITERATIONS.
C                 MDIM  MDIM IS THE USER SUPPLIED DIMENSION OF THE
C                 VECTOR W. IF NMETH=0,MDIM=5*N+2. IF NMETH=1,

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```

C          MDIM=N*(N+7)/2.
C          IDEV  IDEV IS THE USER SUPPLIED NUMBER OF THE OUTPUT
C          DEVICE ON WHICH OUTPUT IS TO BE WRITTEN WHEN
C          IOUTJ0.
C          ACC   ACC IS A USER SUPPLIED ESTIMATE OF MACHINE
C          ACCURACY. A LINEAR SEARCH IS UNSUCCESSFULLY
C          TERMINATED WHEN THE NORM OF THE STEP SIZE
C          BECOMES SMALLER THAN ACC. IN PRACTICE,
C          ACC=10.D-20 HAS PROVED SATISFACTORY. ACC IS
C          DOUBLE PRECISIONED.
C          NMETH NMETH IS THE USER SUPPLIED VARIABLE WHICH
C          CHOOSES THE METHOD OF OPTIMIZATION. IF
C          NMETH=0, A CONJUGATE GRADIENT METHOD IS
C          USED. IF NMETH=1, THE BFGS METHOD IS USED.
C
C REMARKS:  IN ADDITION TO THE SPECIFIED VALUES IN THE ABOVE
C          ARGUMENT LIST, THE USER MUST SUPPLY A SUBROUTINE
C          CALCFG WHICH CALCULATES THE FUNCTION AND GRADIENT AT
C          X AND PLACES THEM IN F AND G(1),...,G(N) RESPECTIVELY.
C          THE SUBROUTINE MUST HAVE THE FORM:
C          SUBROUTINE CALCFG(N,X,F,G)
C          DOUBLE PRECISION X(N),G(N),F
C
C          AN EXAMPLE SUBROUTINE FOR THE ROSENBROCK FUNCTION IS:
C
C          SUBROUTINE CALCFG(N,X,F,G)
C          DOUBLE PRECISION X(N),G(N),F,T1,T2
C          T1=X(2)-X(1)*X(1)
C          T2=1.0-X(1)
C          F=100.0*T1*T1+T2*T2
C          G(1)=-400.0*T1*X(1)-2.0*T2
C          G(2)=200.0*T1
C          RETURN
C          END

```

ALGORITHM 501

Fortran Translation of Algorithm 409, Discrete Chebychev Curve Fit [E2]

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Key Words and Phrases: approximation, polynomial approximation, exchange algorithm, Chebychev approximation, Remez algorithm, minimax polynomial approximation

CR Categories: 5.11, 5.13

Language: Fortran

DESCRIPTION

Subprograms *APPROX* and *EXCH* fit N data pairs stored with independent variable values in array X and dependent variable values in array Y to a minimax polynomial of degree M with parity specified by variable K (0, mixed parity; 1, odd parity; 2, even parity). The polynomial coefficients are returned in ascending order in array A . The magnitude of floating point variable $EPSH$ must be significant when compared with the one on the machine used. The algorithm will fail only if $ABS(EPSH)$ is greater than $HMAX$, the maximum approximation error. Success or failure of a run is indicated in the fixed point variable *EQUAL*. All other arguments are identical with the variables in the original algorithm and must conform to the constraints mentioned there. The variable *NTAPE*, an internal constant in subroutine *APPROX*, specifies the output unit number to which error messages will be directed. To insure the widest possible utility, a primitive level of Fortran was used. By modifying the code, users with access to powerful compilers can reduce both the amount of code and execution time. In particular, the variable *ITEMP* can be eliminated. The algorithm has been used on Sigma 6, Sigma 9, and IBM 360RAX computers.

ACKNOWLEDGMENT

The author is indebted to Herbert Schmitt for suggestions which improved the subprogram.

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ALGORITHM

SUBROUTINE APPROX(M, N, K, X, Y, EPSH, REF, MAXIT, HMAX, H,	APP	10
* A, EQUAL)	APP	20
C THIS SUBROUTINE, IN CONJUNCTION WITH SUBROUTINE EXCH,	APP	30

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C COMPUTES THE MINIMAX (CHEBYCHEV) POLYNOMIAL WHICH FITS	APP 40
C THE DATA IN X AND Y.	APP 50
C THIS SUBPROGRAM IS A TRANSLATION FROM ALGOL OF	APP 60
C H. SCHMITT'S ALGORITHM NUMBER 409 IN CACM, V14,	APP 70
C PP. 355-356 (1971).	APP 80
C THE PRIMARY REFERENCE IS STIEFEL, E. L. NUMERICAL	APP 90
C METHODS OF CHEBYCHEV APPROXIMATION, IN *ON NUMERICAL	APP 100
C APPROXIMATION*, R. LANGER, (EDITOR), UNIVERSITY OF	APP 110
C WISCONSIN PRESS, 1958, PP.217-232.	APP 120
C TRANSLATION BY JOSEPH C. SIMPSON, MARQUETTE UNIVERSITY.	APP 130
C THE EXCHANGE ALGORITHM IS A FINITE ITERATIVE PROCESS	APP 140
C REQUIRING, AT MOST, BICO(N,P) ITERATIONS, WHERE P =	APP 150
C NUMBER OF POINTS FIT IN ONE ITERATION (SEE EXCH FOR THE	APP 160
C VALUE OF P). SINCE BICO(N,P) CAN BE VERY LARGE, IT IS	APP 170
C POSSIBLE THAT THE ROUTINE WILL NOT CONVERGE WITHIN MAXIT	APP 180
C ITERATIONS. THE OTHER POSSIBILITY OF FAILURE OCCURS	APP 190
C WHEN INSUFFICIENT FLOATING POINT PRECISION IS AVAILABLE	APP 200
C FOR THE INPUT DATA CHOSEN.	APP 210
C DESCRIPTION OF INPUT-OUTPUT VARIABLES.	APP 220
C M = DEGREE OF FITTING POLYNOMIAL.	APP 230
C N = NUMBER OF DATA POINTS STORED IN X AND Y.	APP 240
C X = INDEPENDENT VARIABLE VALUES STORED IN ASCENDING ORDER.	APP 250
C Y = DEPENDENT VARIABLE VALUES CORRESPONDING TO ABOVE DATA.	APP 260
C K CONTROLS THE CHARACTER OF THE POLYNOMIAL.	APP 270
C K = 0 IMPLIES MIXED PARITY POLYNOMIAL. X(1) MAY BE ANY	APP 280
C FLOATING POINT VALUE.	APP 290
C K = 1 IMPLIES ODD POLYNOMIAL. X(1) MUST BE .GT.0.	APP 300
C K = 2 IMPLIES EVEN POLYNOMIAL. X(1) MUST BE .GE.0.	APP 310
C EPSH. ABS(EPSH) = TOLERANCE FOR LEVELING. IF EPSH IS	APP 320
C NEGATIVE, REF CONTAINS THE INITIAL POINT SET. ABS(EPSH)	APP 330
C SHOULD BE SIGNIFICANT WHEN COMPARED WITH ONE. A USEFUL	APP 340
C VALUE FOR 24 BIT MANTISSA IS EPSH=2.E-7.	APP 350
C REF = AN ARRAY OF INITIAL POINTS IF EPSH.LT.0. OTHERWISE	APP 360
C IT IS NOT AN INPUT ARRAY. UPON RETURN IT CONTAINS THE	APP 370
C MAXIMUM DEVIATION POINTS.	APP 380
C MAXIT = UPPER LIMIT FOR NUMBER OF EXCHANGE STEPS.	APP 390
C HMAX CONTAINS THE MAXIMUM DEVIATION(OUTPUT).	APP 400
C H CONTAINS THE APPROXIMATION ERRORS(OUTPUT).	APP 410
C A CONTAINS THE POLYNOMIAL COEFFICIENTS IN ORDER OF	APP 420
C INCREASING POWERS(OUTPUT). IF MAXIT IS EXCEEDED, REF	APP 430
C CONTAINS A NEW REFERENCE SET.	APP 440
C DIMENSION REF(REFMAX),DAA(M+1),D(M+2),Z(M+4),A(M+1),	APP 450
C AA(M+1),C(M+2),X(N),Y(N),H(N)	APP 460
C REFMAX(K=0)=M+2. REFMAX(K=1)=INT((M+3)/2).	APP 470
C REFMAX(K=2)=2+M/2.	APP 480
C INPUT VARIABLES ALTERED DURING EXECUTION CAN INCLUDE M,	APP 490
C EPSH,REF.	APP 500
C EQUAL RECORDS THE SUCCESS OR FAILURE OF THE ALGORITHM.	APP 510
C EQUAL=1 SUCCESSFUL.	APP 520
C EQUAL=0 CONVERGENCE OF EXCHANGE PROCESS NOT ACHIEVED.	APP 530
C EQUAL=-1 INPUT ERROR.	APP 540
C EQUAL=-2 ALGORITHM FAILURE. NOTE.. A MAY CONTAIN THE	APP 550
C COLLOCATION POLYNOMIAL.	APP 560
C DESCRIPTION OF SOME LOCAL VARIABLES.	APP 570
C SMALLEST REPRESENTABLE NUMBER. IT IS USED ONLY WHEN	APP 580
C COLLOCATION IS ENCOUNTERED.	APP 590
C THE FOLLOWING VARIABLES ARE DECLARED INTEGER, BUT ARE USED	APP 600
C AS LOGICAL VARIABLES ONLY, .TRUE.=1 AND .FALSE.=0	APP 610
C K0,K1,B1,B2.	APP 620
C Z(I+1) CONTAINS THE ITH POINT OF THE CURRENT REFERENCE	APP 630
C SET. DAA(I+1) CONTAINS THE CORRECTION TO THE COEFFICIENT	APP 640
C OF X**(I).	APP 650
C TO OBTAIN A DOUBLE PRECISION VERSION, CHANGE THE	APP 660
C REAL DECLARATION TO DOUBLE PRECISION, ABS TO DABS, COS	APP 670
C TO DCOS, AND FLOAT TO DFLOAT.	APP 680
C THIS DIMENSION STATEMENT ALLOWS 50 POINTS AND A MIXED	APP 690
C ORDER POLYNOMIAL OF DEGREE 30.	APP 700
DIMENSION XX(50), AA(31), DAA(31), D(32), Z(34), C(32)	APP 710
DIMENSION X(1), Y(1), H(1), REF(1), A(1)	APP 720
DOUBLE PRECISION SD, QD, C, T, D, DT, AA, DAA, XX, MAX	APP 730
REAL EPSH, HMAX, Y, H, A, PI, Q, S, X1, XA, XE, COS	APP 740
INTEGER N, M, K, MAXIT, REF, I, J, P, Q1, Q2, R, K0, K1, B1,	APP 750
* B2, Z, ITEMP, MOLD, HIGH, JJ, II, LOW, EQUAL, NTAPE, B3	APP 760
C MACHINE DEPENDENT CONSTANTS.	APP 770
C NTAPE IS THE PRINTER UNIT NUMBER.	APP 780
NTAPE = 6	APP 790
PI = 3.1415926535	APP 800
MOLD = M	APP 810
C ZERO THE COEFFICIENT ARRAY.	APP 820
IF (K-1) 10, 20, 30	APP 830
10 K0 = 1	APP 840
K1 = 0	APP 850
Q1 = 0	APP 860
Q2 = 1	APP 870
GO TO 50	APP 880
20 K0 = 0	APP 890
K1 = 1	APP 900
Q1 = 1	APP 910
Q2 = 2	APP 920
GO TO 40	APP 930
30 K0 = 0	APP 940


```

      K1 = 0
      Q1 = 0
      Q2 = 2
C ADJUST M ACCORDING TO THE VALUE OF K
      40 M = (M-Q1)/2
      50 P = M + 2
C SCREEN INPUT PARAMETERS.
      IF (P-N) 60, 60, 780
      60 IF (M) 780, 70, 70
      70 IF (K-1) 100, 80, 90
      80 IF (X(1)) 790, 790, 100
      90 IF (X(1)) 790, 100, 100
      100 DO 120 I=2,N
          ITEMP = I - 1
          IF (X(I)-X(ITEMP)) 800, 110, 110
      110 CONTINUE
      120 CONTINUE
          ITEMP = MOLD + 1
          DO 130 I=1,ITEMP
              A(I) = 0.
      130 CONTINUE
          Z(1) = 0
          ITEMP = P + 2
          Z(ITEMP) = N + 1
          IF (EPSH) 140, 170, 170
C IF EPSH.GT.0, BRANCH TO Z SETUP SECTION.
C IF EPSH IS NEGATIVE, TEST REF AND LOAD IT INTO Z.
      140 EPSH = -EPSH
          J = 0
          DO 160 I=1,P
              ITEMP = I + 1
              R = REF(I)
              Z(ITEMP) = R
C BRANCH TO ERROR EXIT UNLESS REF IS MONOTONICALLY
C INCREASING.
          IF (J-R) 150, 810, 810
      150 J = R
      160 CONTINUE
          IF (J-N) 300, 300, 810
C BRANCH AROUND Z SETUP SECTION.
C THIS SECTION LOADS Z WITH THE POINTS CLOSEST TO THE
C CHEBYCHEV ABSCISSAS.
      170 X1 = X(1)
          XE = X(N)
          IF (K0) 190, 190, 180
C IF THE POLYNOMIAL IS NOT MIXED, I.E.HAS A DEFINITE PARITY,
C BRANCH TO 20.
      180 XA = XE + X1
          XE = XE - X1
          Q = PI/FLOAT(M+1)
          GO TO 200
      190 XA = 0.
          XE = XE + XE
          ITEMP = 2*(M+1) + Q1
          Q = PI/FLOAT(ITEMP)
C CALCULATE THE JTH CHEBYCHEV ABSCISSA AND LOAD Z(J+1) WITH
C THE APPROPRIATE INDEX FROM THE DATA ABSCISSAS.
      200 DO 270 JJ=1,P
          J = P + 1 - JJ
          X1 = XA + XE*(COS(Q*FLOAT(P-J)))
          ITEMP = J + 2
          R = Z(ITEMP)
          HIGH = R - 1
          IF (HIGH-2) 230, 210, 210
      210 DO 220 II=2,HIGH
          I = HIGH + 2 - II
          ITEMP = I - 1
          IF (X(I)+X(ITEMP)-X1) 240, 240, 220
C WHEN THE CHEBYCHEV ABSCISSA IS BRACKETED BY TWO INPUT
C ABSCISSAS, BRANCH TO 240.
      220 CONTINUE
      230 I = 1
      240 ITEMP = J + 1
          IF (I-R) 250, 260, 260
      250 Z(ITEMP) = I
          GO TO 270
      260 Z(ITEMP) = R - 1
      270 CONTINUE
C IF THE LOWER CHEBYCHEV ABSCISSAS ARE LESS THAN X(1), LOAD
C THE LOWER ELEMENTS OF Z WITH THE LOWEST POINTS.
          J = 0
      280 J = J + 1
          ITEMP = J + 1
          IF (Z(ITEMP)-J) 290, 300, 300
      290 Z(ITEMP) = J
          GO TO 280
C M1 ENTRY. INITIALIZE VARIABLES TO PREPARE FOR EXCHANGE
C ITERATION.
      300 ITEMP = M + 1
C ZERO THE AA ARRAY.
      DO 310 I=1,ITEMP

```

```

APP 950
APP 960
APP 970
APP 980
APP 990
APP 1000
APP 1010
APP 1020
APP 1030
APP 1040
APP 1050
APP 1060
APP 1070
APP 1080
APP 1090
APP 1100
APP 1110
APP 1120
APP 1130
APP 1140
APP 1150
APP 1160
APP 1170
APP 1180
APP 1190
APP 1200
APP 1210
APP 1220
APP 1230
APP 1240
APP 1250
APP 1260
APP 1270
APP 1280
APP 1290
APP 1300
APP 1310
APP 1320
APP 1330
APP 1340
APP 1350
APP 1360
APP 1370
APP 1380
APP 1390
APP 1400
APP 1410
APP 1420
APP 1430
APP 1440
APP 1450
APP 1460
APP 1470
APP 1480
APP 1490
APP 1500
APP 1510
APP 1520
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APP 1580
APP 1590
APP 1600
APP 1610
APP 1620
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APP 1670
APP 1680
APP 1690
APP 1700
APP 1710
APP 1720
APP 1730
APP 1740
APP 1750
APP 1760
APP 1770
APP 1780
APP 1790
APP 1800
APP 1810
APP 1820
APP 1830
APP 1840
APP 1850

```

```

      AA(I) = 0.
310 CONTINUE
C LOAD H WITH THE ORDINATES AND XX(I) WITH THE ABSCISSAS IF
C THE POLYNOMIAL IS MIXED. IF THE POLYNOMIAL IS EVEN OR ODD
C LOAD XX WITH THE SQUARES OF THE ABSCISSAS.
      DO 340 I=1,N
        H(I) = Y(I)
        Q = X(I)
        IF (K0) 320, 320, 330
320    XX(I) = Q*Q
        GO TO 340
330    XX(I) = Q
340 CONTINUE
      B1 = 0
      B2 = 0
      B3 = 0
      R = -1
      T = 0.
C ITERATION ENTRY. R IS THE ITERATION INDEX.
350 R = R + 1
      S = 1.
C COMPUTATION OF DIVIDED DIFFERENCES SCHEME.
      IF (K1) 380, 380, 360
C IF THE POLYNOMIAL IS ODD, BRANCH TO 380.
360 DO 370 I=1,P
      S = -S
      ITEMP = I + 1
      J = Z(ITEMP)
      Q = X(J)
      C(I) = (H(J)+S*T)/Q
      D(I) = S/Q
370 CONTINUE
      GO TO 400
380 DO 390 I=1,P
      S = -S
      ITEMP = I + 1
      ITEMP = Z(ITEMP)
      C(I) = H(ITEMP) + S*T
      D(I) = S
390 CONTINUE
400 CONTINUE
      DO 420 I=2,P
        DO 410 JJ=I,P
          J = P + I - JJ
          ITEMP = J + 1
          ITEMP = Z(ITEMP)
          QD = XX(ITEMP)
          ITEMP = 2 + J - I
          ITEMP = Z(ITEMP)
          QD = QD - XX(ITEMP)
          ITEMP = J - 1
          C(J) = (C(J)-C(ITEMP))/QD
          D(J) = (D(J)-D(ITEMP))/QD
410 CONTINUE
420 CONTINUE
      DT = -C(P)/D(P)
      T = T + DT
C COMPUTATION OF POLYNOMIAL COEFFICIENTS.
      HIGH = M + 1
      DO 450 II=1,HIGH
        I = HIGH - II
        ITEMP = I + 1
        DAA(ITEMP) = C(ITEMP) + DT*D(ITEMP)
        ITEMP = I + 2
        ITEMP = Z(ITEMP)
        QD = XX(ITEMP)
        LOW = I + 1
        IF (M-LOW) 450, 430, 430
430    DO 440 J=LOW,M
          JJ = J + 1
          DAA(J) = DAA(J) - QD*DAA(JJ)
440 CONTINUE
450 CONTINUE
      DO 460 I=1,HIGH
        AA(I) = AA(I) + DAA(I)
460 CONTINUE
C EVALUATION OF THE POLYNOMIAL TO GET THE APPROXIMATION
C ERRORS.
      MAX = 0.
      DO 540 I=1,N
        SD = AA(HIGH)
        QD = XX(I)
        IF (-M) 470, 490, 490
470    DO 480 JJ=1,M
          J = M - JJ + 1
          SD = SD*QD + AA(J)
480 CONTINUE
490 CONTINUE
        IF (K1) 510, 510, 500
C IF THE POLYNOMIAL IS ODD MULTIPLY SD BY X(I).

```

```

APP 1860
APP 1870
APP 1880
APP 1890
APP 1900
APP 1910
APP 1920
APP 1930
APP 1940
APP 1950
APP 1960
APP 1970
APP 1980
APP 1990
APP 2000
APP 2010
APP 2020
APP 2030
APP 2040
APP 2050
APP 2060
APP 2070
APP 2080
APP 2090
APP 2100
APP 2110
APP 2120
APP 2130
APP 2140
APP 2150
APP 2160
APP 2170
APP 2180
APP 2190
APP 2200
APP 2210
APP 2220
APP 2230
APP 2240
APP 2250
APP 2260
APP 2270
APP 2280
APP 2290
APP 2300
APP 2310
APP 2320
APP 2330
APP 2340
APP 2350
APP 2360
APP 2370
APP 2380
APP 2390
APP 2400
APP 2410
APP 2420
APP 2430
APP 2440
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APP 2460
APP 2470
APP 2480
APP 2490
APP 2500
APP 2510
APP 2520
APP 2530
APP 2540
APP 2550
APP 2560
APP 2570
APP 2580
APP 2590
APP 2600
APP 2610
APP 2620
APP 2630
APP 2640
APP 2650
APP 2660
APP 2670
APP 2680
APP 2690
APP 2700
APP 2710
APP 2720
APP 2730
APP 2740
APP 2750

```

```

500 SD = SD*X(I) APP 2760
510 QD = Y(I) - SD APP 2770
      H(I) = QD APP 2780
      IF (DABS(QD)-MAX) 530, 530, 520 APP 2790
C LOAD MAX WITH THE LARGEST MAGNITUDE OF THE APPROXIMATION APP 2800
C ARRAY. APP 2810
520 MAX = DABS(QD) APP 2820
530 CONTINUE APP 2830
540 CONTINUE APP 2840
C TEST FOR ALTERNATING SIGNS. APP 2850
      ITEMP = Z(2) APP 2860
      IF (H(ITEMP)) 550, 560, 600 APP 2870
550 J = 1 APP 2880
      GO TO 610 APP 2890
560 J = 0 APP 2900
C THIS REPRESENTS A CASE WHERE THE POLYNOMIAL EXACTLY APP 2910
C PREDICTS A DATA POINT. APP 2920
      WRITE (NTAPE,99998) APP 2930
      IF (B3) 570, 570, 690 APP 2940
570 B3 = 1 APP 2950
      IF (EPSH-MAX) 580, 710, 710 APP 2960
580 WRITE (NTAPE,99999) APP 2970
      LOW = (N+1)/2 - (P+1)/2 + 1 APP 2980
      HIGH = LOW + P APP 2990
      DO 590 I=LOW,HIGH APP 3000
          Z(I) = I - 1 APP 3010
590 CONTINUE APP 3020
      GO TO 350 APP 3030
600 J = -1 APP 3040
610 CONTINUE APP 3050
      DO 670 I=2,P APP 3060
          ITEMP = I + 1 APP 3070
          ITEMP = Z(ITEMP) APP 3080
C IF(H(ITEMP)) 339,340,341 APP 3090
          IF (H(ITEMP)) 620, 560, 630 APP 3100
620 JJ = -1 APP 3110
          GO TO 640 APP 3120
630 JJ = 1 APP 3130
640 IF (J-JJ) 650, 660, 650 APP 3140
C ERROR ENTRY. INSUFFICIENT ACCURACY FOR CALCULATION. APP 3150
650 B1 = 1 APP 3160
          GO TO 720 APP 3170
660 J = -J APP 3180
670 CONTINUE APP 3190
C SEARCH FOR ANOTHER REFERENCE. APP 3200
      CALL EXCH(N, P, H, EPSH, Z, EQUAL) APP 3210
      IF (EQUAL) 680, 680, 720 APP 3220
680 IF (R-MAXIT) 350, 700, 700 APP 3230
690 B3 = -1 APP 3240
          GO TO 720 APP 3250
700 B2 = 1 APP 3260
          GO TO 720 APP 3270
710 WRITE (NTAPE,99988) MAX APP 3280
C END OF ITERATION LOOP. APP 3290
C M2 ENTRY. LOAD OUTPUT VARIABLES AND RETURN. APP 3300
720 HIGH = M + 1 APP 3310
C LOAD THE COEFFICIENTS INTO THE A ARRAY. APP 3320
      DO 730 I=1,HIGH APP 3330
          ITEMP = Q1 + Q2*(I-1) + 1 APP 3340
          A(ITEMP) = AA(I) APP 3350
730 CONTINUE APP 3360
C LOAD REF WITH THE FINAL REFERENCE POINTS. APP 3370
      DO 740 I=1,P APP 3380
          ITEMP = I + 1 APP 3390
          REF(I) = Z(ITEMP) APP 3400
740 CONTINUE APP 3410
      HMAX = MAX APP 3420
      MAXIT = R APP 3430
      IF (B3) 840, 750, 750 APP 3440
750 IF (B1) 760, 760, 820 APP 3450
760 IF (B2) 770, 770, 830 APP 3460
770 M = MOLD APP 3470
C NORMAL EXIT APP 3480
      RETURN APP 3490
C ERROR EXITS. APP 3500
780 WRITE (NTAPE,99990) APP 3510
      WRITE (NTAPE,99996) N, MOLD, K APP 3520
      EQUAL = -1 APP 3530
      GO TO 770 APP 3540
790 WRITE (NTAPE,99990) APP 3550
      WRITE (NTAPE,99995) K, X(1) APP 3560
      EQUAL = -1 APP 3570
      GO TO 770 APP 3580
800 WRITE (NTAPE,99990) APP 3590
      WRITE (NTAPE,99994) (X(I),I=1,N) APP 3600
      EQUAL = -1 APP 3610
      GO TO 770 APP 3620
810 WRITE (NTAPE,99990) APP 3630
      WRITE (NTAPE,99993) (REF(I),I=1,N) APP 3640
      EQUAL = -1 APP 3650
      GO TO 770 APP 3660

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820 WRITE (NTAPE,99997)	APP 3670
WRITE (NTAPE,99992) HMAX	APP 3680
EQUAL = -2	APP 3690
GO TO 770	APP 3700
830 WRITE (NTAPE,99997)	APP 3710
ITEMP = M + 2	APP 3720
WRITE (NTAPE,99991) MAXIT, (REF(I),I=1,ITEMP)	APP 3730
EQUAL = 0	APP 3740
GO TO 770	APP 3750
840 WRITE (NTAPE,99997)	APP 3760
WRITE (NTAPE,99989)	APP 3770
EQUAL = -2	APP 3780
GO TO 770	APP 3790
99999 FORMAT (1H+, 26X, 35H ONE MORE ATTEMPT WILL BE MADE USING,	APP 3800
* 19H THE MIDDLE POINTS.)	APP 3810
99998 FORMAT (25H COLLOCATION HAS OCCURRED.)	APP 3820
99997 FORMAT (26H ALGORITHM APPROX FAILURE.)	APP 3830
99996 FORMAT (10X, 28H INSUFFICIENT INFORMATION. N=, I3, 3H M=, I3,	APP 3840
* 3H K=, I3)	APP 3850
99995 FORMAT (10X, 37HPOLYNOMIAL PARITY - ABSCISSA CONFLICT.,	APP 3860
* 6H K =, I3, 7H X(1)=, F20.10)	APP 3870
99994 FORMAT (10X, 40H ABSCISSAS OUT OF ORDER. X ARRAY EQUALS	APP 3880
* /(7E17.7))	APP 3890
99993 FORMAT (10X, 40H INITIAL POINT ARRAY NOT MONOTONIC INCREA,	APP 3900
* 23HSING. REF ARRAY EQUALS/(1319))	APP 3910
99992 FORMAT (1X, 39H APPROXIMATION ERRORS AT POINTS OF REFER,	APP 3920
* 30HENCE DO NOT ALTERNATE IN SIGN./23H SUSPECT INSUFFICIENT W,	APP 3930
* 11H ORD LENGTH., 16H MAXIMUM ERROR=, E15.7/14H IF POLYNOMIAL,	APP 3940
* 34H WAS ASSUMED TO BE OF MIXED PARITY, 18H, CHECK FOR EVEN O,	APP 3950
* 13HR ODD PARITY.)	APP 3960
99991 FORMAT (31H MAXIMUM NUMBER OF ITERATIONS, , I6, 10H THE REFER,	APP 3970
* 23HENCE ARRAY OBTAINED IS /(1319))	APP 3980
99990 FORMAT (31H INPUT ERROR SUBPROGRAM APPROX.)	APP 3990
99989 FORMAT (1H+, 26X, 18H APPROX TERMINATES.)	APP 4000
99988 FORMAT (1H+, 26X, 6H MAX=, E15.7, 23H .LT.HMIN. NORMAL EXIT,	APP 4010
* 1H.)	APP 4020
END	APP 4030
SUBROUTINE EXCH(N, P, H, EPSH, Z, EQUAL)	EXC 10
C DIMENSION H(N), Z(P+2)	EXC 20
C N = NUMBER OF POINTS.	EXC 30
C P = NUMBER OF REFERENCE POINTS.	EXC 40
C EPSH = APPROXIMATION ERROR STANDARD.	EXC 50
C Z= INPUT OLD REFERENCE INDICES.	EXC 60
C OUTPUT NEW REFERENCE INDICES.	EXC 70
C Z(1) SHOULD CONTAIN ZERO.	EXC 80
C Z(P+2) SHOULD CONTAIN (N+1).	EXC 90
C Z(I+1) CONTAINS THE ITH REFERENCE POINT INDEX.	EXC 100
C H = ARRAY OF APPROXIMATION ERRORS.	EXC 110
C EQUAL = 0 IMPLIES NORMAL EXCHANGE.	EXC 120
C EQUAL = 1 IMPLIES OLD AND NEW REFERENCE SETS ARE EQUAL.	EXC 130
C THE APPROXIMATION ERRORS ARE COMPARED RELATIVE TO EPSH.	EXC 140
C DESCRIPTION OF SOME LOCAL VARIABLES.	EXC 150
C HZ1 CONTAINS THE LOW POINT APPROXIMATION ERROR.	EXC 160
C HZP CONTAINS THE HIGH POINT APPROXIMATION ERROR.	EXC 170
C MAXR CONTAINS THE LARGEST APPROXIMATION ERROR ABOVE THE	EXC 180
C REFERENCE POINT SET. INDR IS THE POINT INDEX FOR MAXR.	EXC 190
C MAXL CONTAINS THE LARGEST APPROXIMATION ERROR BELOW THE	EXC 200
C REFERENCE POINT SET. INDL IS THE POINT INDEX FOR MAXL.	EXC 210
C ITEMP IS A FIXED POINT VARIABLE USED FOR TEMPORARY STORAGE	EXC 220
C CODING FOR HIGHER LEVEL FORTRAN COULD ELIMINATE ITEMP.	EXC 230
C TO OBTAIN A DOUBLE PRECISION VERSION, CHANGE THE REAL	EXC 240
C DECLARATION TO DOUBLE PRECISION AND ABS TO DABS.	EXC 250
DIMENSION Z(1), H(1)	EXC 260
INTEGER Z, N, P, I, J, L, INDEX, INDL, INDR, ZE, ITEMP, LOW,	EXC 270
* HIGH, EQUAL, II	EXC 280
REAL EPSH, H, HZ1, HZP, MAX, MAXL, MAXR, SIG	EXC 290
EQUAL = 0	EXC 300
L = 0	EXC 310
ITEMP = Z(2)	EXC 320
C ARBITRARILY CHOSEN EQUAL TO THE SIGN OF THE INPUT POINT.	EXC 330
C THIS WILL BE ADJUSTED LATER IF NECESSARY.	EXC 340
IF (H(ITEMP)) 10, 10, 20	EXC 350
10 SIG = 1.	EXC 360
GO TO 30	EXC 370
20 SIG = -1.	EXC 380
30 DO 90 I=1,P	EXC 390
C DO 90 PRESCANS Z TO INSURE IT IS A PROPER CHOICE, I. E.	EXC 400
C RESETS Z IF NECESSARY SO THAT MAXIMUM ERROR POINTS ARE	EXC 410
C CHOSEN, GIVEN THE SIGN CONVENTION MENTIONED ABOVE.	EXC 420
C IN ORDER TO WORK PROPERLY, THIS SECTION REQUIRES Z(1)=0	EXC 430
C AND Z(P+2)=N+1.	EXC 440
MAX = 0.	EXC 450
SIG = -SIG	EXC 460
ITEMP = I + 2	EXC 470
ZE = Z(ITEMP) - 1	EXC 480
LOW = Z(I) + 1	EXC 490
C SCAN THE OPEN POINT INTERVAL CONTAINING ONLY THE ITH	EXC 500
C INITIAL REFERENCE POINT. IN THE INTERVAL PICK THE	EXC 510


```

C IF A POINT OUTSIDE THE PRESENT REFERENCE SET MUST BE          EXC 1420
C INCLUDED, (I.E. THE SIGN OF THE FIRST POINT, ASSUMED IN      EXC 1430
C THE DO 90 SECTION IS WRONG) SHIFT TO THE SIDE OF LARGEST    EXC 1440
C RELATIVE ERROR FIRST.                                       EXC 1450
C CHECK THE OTHER SIDE.                                       EXC 1460
  320 IF (MAXL-MAXR) 350, 350, 330                               EXC 1470
  330 IF (MAXL-HZP) 340, 340, 420                               EXC 1480
  340 IF (MAXR-HZ1) 310, 370, 370                               EXC 1490
  350 IF (MAXR-HZ1) 360, 360, 370                               EXC 1500
  360 IF (MAXL-HZP) 310, 420, 420                               EXC 1510
C SHR ENTRY.                                                  EXC 1520
C THIS SECTION INSERTS A POINT FROM ABOVE THE PRESCAN        EXC 1530
C POINT SET.                                                  EXC 1540
  370 INDEX = Z(2)                                             EXC 1550
      DO 380 I=2,P                                             EXC 1560
C SHIFT POINT SET DOWN, DROPPING THE LOWEST POINT.          EXC 1570
      ITEMP = I + 1                                           EXC 1580
      Z(I) = Z(ITEMP)                                         EXC 1590
  380 CONTINUE                                               EXC 1600
      ITEMP = P + 1                                           EXC 1610
C ADD THE NEW HIGH POINT.                                     EXC 1620
      Z(ITEMP) = INDR                                         EXC 1630
      IF (MAXL) 310, 310, 390                                  EXC 1640
C IF MAXL=0, RETURN.                                         EXC 1650
C IF MAXL.GT.0 REPLACE REFERENCE POINTS FROM THE LEFT,      EXC 1660
C STOPPING THE FIRST TIME THE CANDIDATE FOR REPLACEMENT IS  EXC 1670
C GREATER IN MAGNITUDE THAN THE PROSPECTIVE REPLACEE.      EXC 1680
C ALTERNATION OF SIGNS IS PRESERVED BECAUSE NON-REPLACEMENT EXC 1690
C IMMEDIATELY TERMINATES THE PROCESS.                       EXC 1700
  390 DO 410 I=2,P                                           EXC 1710
      ITEMP = Z(I)                                           EXC 1720
      IF (ABS(H(INDL))-ABS(H(ITEMP))) 310, 400, 400          EXC 1730
  400  J = Z(I)                                               EXC 1740
      Z(I) = INDL                                             EXC 1750
      INDL = INDEX                                           EXC 1760
      INDEX = J                                               EXC 1770
  410 CONTINUE                                               EXC 1780
      GO TO 310                                              EXC 1790
C ENTRY SHL. THIS SECTION INSERTS A POINT FROM BELOW THE   EXC 1800
C PRESCAN POINT SET.                                         EXC 1810
  420 ITEMP = P + 1                                           EXC 1820
      INDEX = Z(ITEMP)                                        EXC 1830
C SHIFT REFERENCE POINT SET UP BY ONE.                       EXC 1840
      DO 430 II=2,P                                          EXC 1850
          I = P + 2 - II                                     EXC 1860
          ITEMP = I + 1                                     EXC 1870
  430 CONTINUE                                               EXC 1890
C STORE LOWEST POINT IN Z(2)                                  EXC 1900
      Z(2) = INDL                                           EXC 1910
      IF (MAXR) 310, 310, 440                                  EXC 1920
C IF MAXR=0, RETURN.                                         EXC 1930
C IF MAXR.GT.0, START REPLACING REFERENCE POINTS FROM RIGHT, EXC 1940
C STOPPING THE FIRST TIME THE CANDIDATE FOR REPLACEMENT    EXC 1950
C IS GREATER IN MAGNITUDE THAN THE PROSPECTIVE REPLACEE.   EXC 1960
  440 DO 460 II=2,P                                          EXC 1970
      I = P + 2 - II                                       EXC 1980
      ITEMP = I + 1                                         EXC 1990
      HIGH = Z(ITEMP)                                       EXC 2000
      IF (ABS(H(INDR))-ABS(H(HIGH))) 310, 450, 450          EXC 2010
  450  J = Z(ITEMP)                                         EXC 2020
      Z(ITEMP) = INDR                                       EXC 2030
      INDR = INDEX                                         EXC 2040
      INDEX = J                                             EXC 2050
  460 CONTINUE                                               EXC 2060
      GO TO 310                                              EXC 2070
      END                                                    EXC 2080

```

ACM Transactions on Mathematical Software, Vol. 4, No. 1, March 1978, Page 95.

REMARK ON ALGORITHM 501

Fortran Translation of Algorithm 409. Discrete Chebychev Curve Fit [E2]
 [Joseph C. Simpson, *ACM Trans. Math. Software* 2, 1 (March 1976), 95-97]

R. Futrell [Recd 16 February 1977]

Harris Corporation, P.O. Box 37, Melbourne, FL 32901

The following should be inserted just above the statement labeled 430 in subroutine
*EXCH*¹:

Z(ITEMP) = Z(I) EXC 1880

With this correction this algorithm has been used to fit first, second, and third
 degree curves to experimental data on the following machines: Harris SLASH 5,
 Honeywell 600 series, and Univac 1108.

¹Subroutine *EXCH* is published in the "Collected Algorithms from ACM"; to insert line
 EXC 1880, see page 501-P8-R1

ALGORITHM 502

Dependence of Solution of Nonlinear Systems on a Parameter [C5]

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Key Words and Phrases: solution of nonlinear equation, differentiation with respect to a parameter, one-parameter imbedding, dependence on parameter

CR Categories: 5.15

Language: Fortran

DESCRIPTION

The Fortran algorithm *DERPAR* is a subroutine suitable for evaluation of the dependence of the solution x_1, \dots, x_n of a system of equations

$$\begin{aligned} f_1(x_1, \dots, x_n, \alpha) &= 0 \\ &\vdots \\ f_n(x_1, \dots, x_n, \alpha) &= 0 \end{aligned} \quad (1)$$

on the parameter α . The modified method of Daviděnko [1, 2, 3], which applies the implicit function theorem, is used in combination with the Newton method and Adams integration formulas. A similar procedure was used by Deist and Sefor [4], who made use of an introduced parameter.

To establish the dependence $x(\alpha)$, the following set of differential equations can be derived:

$$\frac{dx}{d\alpha} = -\Gamma^{-1}(x, \alpha) \frac{\partial f}{\partial \alpha}, \quad x(\alpha_0) = x_0. \quad (2)$$

Here

$$\Gamma = \left\{ \frac{\partial f_i}{\partial x_j} \right\}, \quad \frac{\partial f}{\partial \alpha} = \left(\frac{\partial f_1}{\partial \alpha}, \dots, \frac{\partial f_n}{\partial \alpha} \right)^T, \quad f(x_0, \alpha_0) = 0.$$

If $\Gamma(x(\alpha), \alpha)$ is a regular matrix for $\langle \alpha_0, \alpha_1 \rangle$, then $x(\alpha)$ obtained by integration of (2) satisfies

$$f(x(\alpha), \alpha) = 0, \quad \alpha \in \langle \alpha_0, \alpha_1 \rangle. \quad (3)$$

However, at a branching point the matrix Γ is singular and the integration procedure mentioned above fails.

Therefore, the method must be adapted to be able to handle branching points, i.e. to evaluate the whole dependence $x(\alpha)$, which is a continuously smooth curve in the $(n + 1)$ -th dimensional space $(x - \alpha)$. A parameterization with respect to

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ACM Transactions on Mathematical Software, Vol. 2, No. 1, March 1976, Pages 98-107.

the arc length of the solution locus seems to be a powerful technique which overcomes the difficulties [5].

Differentiation of eq. (1) with respect to t gives

$$\frac{df_i}{dt} = \sum_{j=1}^n \frac{\partial f_i}{\partial x_j} \frac{dx_j}{dt} + \frac{\partial f_i}{\partial \alpha} \frac{d\alpha}{dt} = 0, \quad i = 1, 2, \dots, n. \quad (4)$$

An additional equation

$$\left(\frac{dx_1}{dt}\right)^2 + \dots + \left(\frac{dx_n}{dt}\right)^2 + \left(\frac{d\alpha}{dt}\right)^2 = 1 \quad (5)$$

determines the parameter t as the arc length of the curve $x(\alpha)$ in the space $(x - \alpha)$. The initial conditions for eqs. (4) and (5) are in the form

$$t = 0 : x = x_0, \alpha = \alpha_0. \quad (6)$$

Equation (4) forms a system of n linear equations for $n+1$ unknowns dx_i/dt , $i = 1, \dots, n$, $d\alpha/dt$, for each t . Suppose that the matrix

$$\Gamma_k = \begin{pmatrix} \frac{\partial f_1}{\partial x_1}, \dots, \frac{\partial f_1}{\partial x_{k-1}}, \frac{\partial f_1}{\partial x_{k+1}}, \dots, \frac{\partial f_1}{\partial x_{n+1}} \\ \frac{\partial f_2}{\partial x_1}, \dots \\ \vdots \\ \frac{\partial f_n}{\partial x_1}, \dots, \frac{\partial f_n}{\partial x_{k-1}}, \frac{\partial f_n}{\partial x_{k+1}}, \dots, \frac{\partial f_n}{\partial x_{n+1}} \end{pmatrix} \quad (7)$$

is regular (for certain values of t and k , $1 \leq k \leq n+1$). For clarity here, we have denoted

$$x_{n+1} = \alpha. \quad (8)$$

Equation (4) can then be presolved with respect to the unknowns $dx_1/dt, \dots, dx_{k-1}/dt, dx_{k+1}/dt, \dots, dx_{n+1}/dt$, depending on dx_k/dt in the form

$$\frac{dx_i}{dt} = \beta_i \frac{dx_k}{dt}, \quad i = 1, 2, \dots, k-1, k+1, \dots, n+1. \quad (9)$$

In the program the unknowns with a subscript k (dx_k/dt and x_k) are called the independent variables. The value of the subscript k is determined in the Gaussian elimination procedure using controlled pivoting. The result depends on the values of $PREF(I)$.

On substituting (9) into (5) we obtain

$$\left(\frac{dx_k}{dt}\right)^2 = \left(1 + \sum_{i=1, i \neq k}^{n+1} \beta_i^2\right)^{-1}. \quad (10)$$

$\gamma = 1000$, $B = 22$, $\beta_1 = \beta_2 = 2$, $\theta_{c1} = \theta_{c2} = 0$; $x_0 = (0;0;0;0)$, $\alpha_0 = 0$,
 $PREF(I) = (1, 0.1, 1, 1, 0.02)$, $HH = 0.02$, $EPS = 10^{-4}$, $MXADMS = 4$,
 $NCORR = 4$, $NCRAD = 0$, $NDIR(I) = (1, 1, 1, 1, 1)$.

Calculated points on the curve $x_2(\alpha)$ are numbered. For the points 1-12, 61-70, 114-140, 155-... , α was independent ($k = 5$); otherwise x_2 ($k = 2$) plays the role of an independent variable. Altogether 332 evaluations of the subroutine *FCTN* for calculation of 162 points was necessary. However, by choosing $NCORR = 0$ and $NCRAD = 1$, the number of evaluations of *FCTN* was lowered to 163, and in fact the same curve was obtained.

The sign of dx_k/dt is given by the orientation of the parameter t along the curve. All derivatives dx_i/dt are then determined by (9).

The Adams-Bashforth explicit multistep method with an automatic change in the order of approximation is used for the integration of differential equations (9) and (10). It is possible to use some other technique, e.g. the Runge-Kutta method. However, it seems that our technique is more suitable when considering the time-

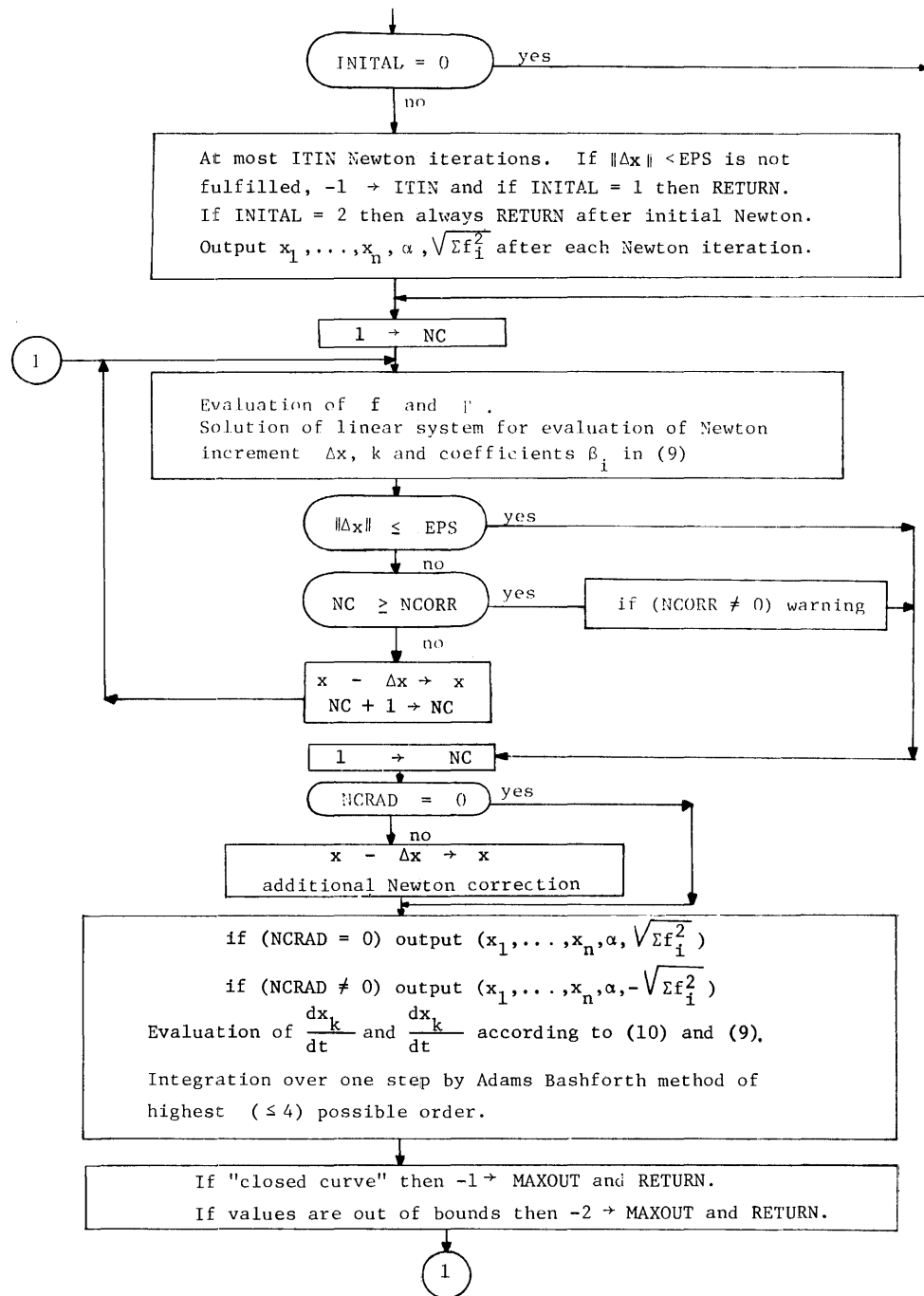


Fig. 1. Schematic flowchart of method

consuming evaluation of the Jacobian matrix and the matrix inversion procedure in each step.

In the course of integration the truncation error causes a deviation between the calculated solution $x(t)$ and the correct profiles $x(t)$. The Newton method for variables $x' = (x_1, x_2, \dots, x_{k-1}, x_{k+1}, \dots, x_{n+1})^T$ is therefore used to improve the calculated profiles:

$$x'_{\text{new}} - x'_{\text{old}} = -\Gamma_k^{-1}f.$$

The higher accuracy of numerical integration lowers the computer time expenditure only because of the low number of Newton iterations.

The algorithm is described by means of the schematic flow diagram given in

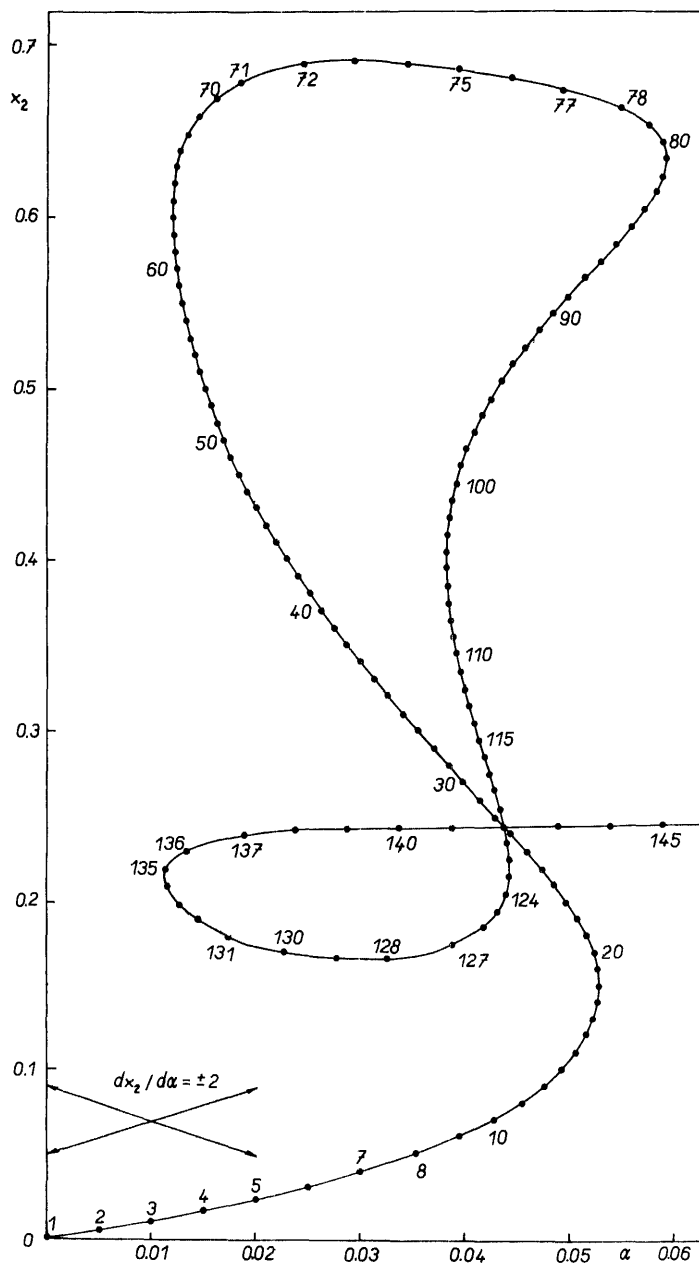
Fig. 2. The dependence $x_2(\alpha)$

Figure 1. The specification of the individual dummy (formal) parameters of the subroutine *DERPAR* is given in the subroutine heading.

The particular form of eq. (1) has to be programmed by the user as a subroutine *FCTN*. To illustrate the construction of the subroutine *FCTN*, an example is presented for a set of four nonlinear equations:

$$f_1 = \alpha(1 - x_3) \exp(10x_1/(1 + 10x_1/\gamma)) - x_3 = 0$$

$$f_2 = \alpha B(1 - x_3) \exp(10x_1/(1 + 10x_1/\gamma)) + \beta_1 \theta_{c1} - 10(1 + \beta_1)x_1 = 0$$

$$f_3 = x_3 - x_4 + \alpha(1 - x_4) \exp(10x_2/(1 + 10x_2/\gamma)) = 0$$

$$f_4 = 10x_1 - 10(1 + \beta_2)x_2 + \alpha B(1 - x_4) \exp(10x_2/(1 + 10x_2/\gamma)) + \beta_2 \theta_{c2} = 0.$$

Here γ , B , β_1 , β_2 , θ_{c1} , θ_{c2} , and α are physical parameters. The dependence $x_2(\alpha)$ is shown in Figure 2. Generally speaking, it is always convenient to transform the unknowns $x_1, x_2, \dots, x_n, \alpha$ so that their values are of similar magnitude.

Applicability. This program can be used in the computation of the dependence of a solution on a parameter for sets of nonlinear equations, for nonlinear boundary value problems for ordinary differential equations [6, 7], for difference equations [8], etc.

ACKNOWLEDGMENTS

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ALGORITHM

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SUBROUTINE DERPAR(N, X, XLOW, XUPP, EPS, W, INITAL, ITIN, HH, DER 10
* HMAX, PREF, NDIR, E, MXADMS, N CORR, NCRAD, NOUT, OUT, DER 20
* MAXOUT, NPRNT) DER 30
DIMENSION X(11), XLOW(11), XUPP(11), W(11), HMAX(11), DER 40
* PREF(11), NDIR(11), OUT(10,12), F(11), G(10,11), BETA(11), DER 50
* MARK(11), DXDT(11) DER 60
C OBTAINING OF DEPENDENCE OF SOLUTION X(ALFA) OF EQUATION DER 70
C F ( X , ALFA ) = 0 DER 80
C ON PARAMETER ALFA BY MODIFIED METHOD OF DIFFERENTIATION DER 90
C WITH RESPECT TO PARAMETER DER 100
C N - NUMBER OF UNKNOWNNS X(I) DER 110
C X(1),...,X(N) - INITIAL VALUES OF X(I), AFTER RETURN FINAL VALUES DER 120
C OF X(I) DER 130
C X(N+1) - INITIAL VALUE OF PARAMETER ALFA, AFTER RETURN FINAL VALUE DER 140
C OF ALFA DER 150
C XLOW(1),...,XLOW(N) - LOWER BOUNDS FOR X(I) DER 160
C XUPP(1),...,XUPP(N) - UPPER BOUNDS FOR X(I) DER 170
C XLOW(N+1),XUPP(N+1) - LOWER AND UPPER BOUNDS FOR ALFA DER 180
C IF XLOW OR XUPP IS EXCEEDED, THEN END OF DERPAR DER 190
C AND MAXOUT=-2 AFTER RETURN DER 200
C EPS - ACCURACY DESIRED IN NEWTON ITERATION FOR DER 210
C SUM OF (W(I)*ABS(XNEW(I)-XOLD(I))), I=1,...,N+1 DER 220
C W(1),...,W(N+1) - WEIGHTS USED IN TERMINATION CRITERION OF NEWTON DER 230
C PROCESS DER 240
C INITAL - IF (INITAL.NE.0) THEN SEVERAL STEPS IN NEWTON ITERATION DER 250
C ARE MADE BEFORE COMPUTATION IN ORDER TO INCREASE DER 260
C ACCURACY OF INITIAL POINT- DER 270
C IF (INITAL.EQ.1.AND.EPS-ACCURACY IS NOT FULFILLED IN ITIN DER 280
C ITERATIONS) THEN RETURN. IF (INITAL.EQ.2) THEN ALWAYS RETURN DER 290
C AFTER INITIAL NEWTON ITERATION, RESULTS ARE IN X. DER 300
C IF (INITAL.EQ.3) THEN CONTINUE IN DERPAR AFTER INITIAL NEWTON. DER 310
C IF (INITAL.EQ.0) THEN NO INITIAL NEWTON ITERATION IS USED. DER 320
C ITIN - MAXIMAL NUMBER OF INITIAL NEWTON ITERATIONS. IF DER 330
C EPS-ACCURACY IS NOT FULFILLED IN ITIN ITERATIONS THEN DER 340
C ITIN=-1 AFTER RETURN. DER 350
C HH - INTEGRATION STEP ALONG ARC LENGTH OF SOLUTION LOCUS DER 360
C HMAX(1),...,HMAX(N+1) - UPPER BOUNDS FOR INCREMENTS OF X(I) IN DER 370
C ONE INTEGRATION STEP (APPROXIMATION ONLY) DER 380
C PREF(1),...,PREF(N+1) - PREFERENCE NUMBERS (EXPLANATION SEE IN DER 390
C SUBR.GAUSE) DER 400
C NDIR(1),...,NDIR(N+1) - INITIAL CHANGE OF X(I) IS POSITIVE ALONG DER 410
C SOLUTION LOCUS (CURVE) IF NDIR(I)=1 AND NEGATIVE IF DER 420
C NDIR(I)=-1. DER 430
C E - CRITERION FOR TEST ON CLOSED CURVE, IF DER 440
C (SUM OF (W(I)*ABS(X(I)-XINITIAL(I))),I=1,...,N+1).LE.E) DER 450
C THEN CLOSED CURVE MAY BE EXPECTED. DER 460
C MXADMS - MAXIMAL ORDER OF ADAMS-BASHFORTH FORMULA, DER 470
C 1.LE.MXADMS.LE.4. DER 480

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C NCORR - MAXIMAL NUMBER OF NEWTON CORRECTIONS AFTER PREDICTION          DER 490
C   BY ADAMS-BASHFORTH METHOD.                                           DER 500
C NCRAD - IF (NCRAD.NE.0) THEN ADDITIONAL NEWTON CORRECTION             DER 510
C   WITHOUT NEW COMPUTATION OF JACOBIAN MATRIX IS USED.                 DER 520
C NOUT - AFTER RETURN NUMBER OF CALCULATED POINTS ON THE CURVE         DER 530
C   X(ALFA), NOUT.LE.MAXOUT.                                            DER 540
C OUT(J,1),OUT(J,2),...,OUT(J,N+1) - J-TH POINT X(1),...,X(N),ALFA    DER 550
C   ON CURVE X(ALFA)                                                    DER 560
C OUT(J,N+2) - VALUE OF SQRT(SUM OF SQUARES OF F).                      DER 570
C IF (OUT(J,N+2).LT.0.0) THEN ABS(OUT(J,N+2)) CORRESPONDS TO X         DER 580
C AND ALFA NOT EXACTLY, BECAUSE ADDITIONAL NEWTON CORRECTION WAS      DER 590
C USED (NCRAD.NE.0).                                                    DER 600
C VALUES F(I) ARE NOT COMPUTED FOR X AND ALFA PRINTED AND/OR         DER 610
C STORED AND THEREFORE LAST TIME COMPUTED VALUE OF SQRT(SUM OF      DER 620
C SQUARES OF F) IS ON OUR DISPOSAL ONLY.                                DER 630
C MAXOUT - MAXIMAL NUMBER OF CALCULATED POINTS ON CURVE X(ALFA).       DER 640
C IF MAXOUT AFTER RETURN EQUALS TO-                                     DER 650
C   -1 - THEN CLOSED CURVE X(ALFA) MAY BE EXPECTED                     DER 660
C   -2 - THEN BOUND XLOW OR XUPP WAS EXCEEDED                          DER 670
C   -3 - THEN SINGULAR JACOBIAN MATRIX OCCURRED, ITS RANK IS         DER 680
C   .LT. N.                                                             DER 690
C NPRNT - IF (NPRNT.EQ.3) THEN RESULTING POINTS ON CURVE X(ALFA)      DER 700
C ARE IN ARRAY OUT( , ) AFTER RETURN. IF (NPRNT.EQ.1) THEN THESE     DER 710
C POINTS ARE PRINTED ONLY. IF (NPRNT.EQ.2) THEN THESE POINTS ARE     DER 720
C BOTH PRINTED AND IN ARRAY OUT.                                       DER 730
C SUBROUTINE FCTN MUST BE PROGRAMMED IN FOLLOWING FORM -              DER 740
C SUBROUTINE FCTN(N,X,F,G)                                             DER 750
C DIMENSION X(11),F(10),G(10,11)                                       DER 760
C F(I)= FI (X(1),X(2),...,X(N),ALFA) FOR I=1,...,N                    DER 770
C G(I,J)= D FI (X(1),...,X(N),ALFA)/ D X(J) FOR I,J=1,...,N         DER 780
C G(I,N+1)= D FI (X(1),...,X(N),ALFA)/ D ALFA FOR I=1,...,N        DER 790
C RETURN                                                                DER 800
C END                                                                    DER 810
C   DATA INDIC, INDSP /1H*,1H /                                       DER 820
C LW IS PRINTER DEVICE NUMBER                                         DER 830
C   N1 = N + 1                                                         DER 840
C   LW = 3                                                             DER 850
C   IF (INITAL) 10, 60, 10                                           DER 860
C INITIAL NEWTON ITERATIONS                                           DER 870
C 10 DO 40 L=1,ITIN                                                  DER 880
C   CALL FCTN(N, X, F, G)                                             DER 890
C   SQUAR = 0.0                                                       DER 900
C   DO 20 I=1,N                                                       DER 910
C     SQUAR = SQUAR + F(I)**2                                         DER 920
C 20 CONTINUE                                                         DER 930
C   LL = L - 1                                                         DER 940
C   SQUAR = SQRT(SQUAR)                                               DER 950
C   IF (NPRNT.NE.3) WRITE (LW,99999) LL, (X(I),I=1,N1), SQUAR     DER 960
C   CALL GAUSE(N, G, F, M, 10, 11, PREF, BETA, K)                   DER 970
C   IF (M.EQ.0) GO TO 310                                             DER 980
C   P = 0.0                                                            DER 990
C   DO 30 J=1,N1                                                      DER 1000
C     X(J) = X(J) - F(J)                                             DER 1010
C     P = P + ABS(F(J))*W(J)                                         DER 1020
C 30 CONTINUE                                                         DER 1030
C   IF (P.LE.EPS) GO TO 50                                           DER 1040
C 40 CONTINUE                                                         DER 1050
C   WRITE (LW,99998) ITIN                                           DER 1060
C   ITIN = -1                                                         DER 1070
C   IF (INITAL.EQ.1) RETURN                                           DER 1080
C 50 IF (NPRNT.NE.3) WRITE (LW,99997) (X(I),I=1,N1)                 DER 1090
C   IF (INITAL.EQ.2) RETURN                                           DER 1100
C AFTER INITIAL NEWTON ITERATIONS                                     DER 1110
C 60 IF (NPRNT.NE.3) WRITE (LW,99996)                                DER 1120
C   KOUT = 0                                                         DER 1130
C   NOUT = 0                                                         DER 1140
C   MADMS = 0                                                         DER 1150
C   NC = 1                                                            DER 1160
C   K1 = 0                                                            DER 1170
C 70 CALL FCTN(N, X, F, G)                                             DER 1180
C   SQUAR = 0.0                                                       DER 1190
C   DO 80 I=1,N                                                       DER 1200
C     SQUAR = SQUAR + F(I)**2                                         DER 1210
C 80 CONTINUE                                                         DER 1220
C   CALL GAUSE(N, G, F, M, 10, 11, PREF, BETA, K)                   DER 1230
C   IF (M.EQ.0) GO TO 310                                             DER 1240
C   IF (K1.EQ.K) GO TO 90                                             DER 1250
C CHANGE OF INDEPENDENT VARIABLE (ITS INDEX = K NOW)                 DER 1260
C   MADMS = 0                                                         DER 1270
C   K1 = K                                                            DER 1280
C 90 SQUAR = SQRT(SQUAR)                                             DER 1290
C   IF (NCRAD.EQ.1) SQUAR = -SQUAR                                    DER 1300
C   P = 0.0                                                            DER 1310
C   DO 100 I=1,N1                                                     DER 1320
C     P = P + W(I)*ABS(F(I))                                         DER 1330
C 100 CONTINUE                                                       DER 1340
C   IF (P.LE.EPS) GO TO 130                                           DER 1350
C   IF (NC.GE.NCORR) GO TO 120                                       DER 1360
C   DO 110 I=1,N1                                                     DER 1370
C     X(I) = X(I) - F(I)                                             DER 1380
C 110 CONTINUE                                                       DER 1390
C ONE ITERATION IN NEWTON METHOD                                       DER 1400

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      NC = NC + 1
      GO TO 70
120  IF (NCORR.EQ.0) GO TO 130
      WRITE (LW,99995) NCORR, P
130  NC = 1
      IF (NCRAD.EQ.0) GO TO 150
C ADDITIONAL NEWTON CORRECTION
      DO 140 I=1,N1
        X(I) = X(I) - F(I)
140  CONTINUE
150  NOUT = NOUT + 1
      DO 160 I=1,N1
        MARK(I) = INDSP
160  CONTINUE
      MARK(K) = INDIC
      IF (NPRNT.EQ.3) GO TO 170
      WRITE (LW,99994)
      WRITE (LW,99993) (X(I),MARK(I),I=1,N1), SQUAR
170  IF (NPRNT.EQ.1) GO TO 200
      IF (NOUT.LE.100) GO TO 180
      WRITE (LW,99992)
      RETURN
180  DO 190 I=1,N1
        OUT(NOUT,I) = X(I)
190  CONTINUE
      OUT(NOUT,N+2) = SQUAR
      GO TO 210
200  IF (NOUT.EQ.1) GO TO 180
210  IF (NOUT.GE.MAXOUT) RETURN
      DO 220 I=1,N1
        IF (X(I).LT.XLOW(I) .OR. X(I).GT.XUPP(I)) GO TO 300
220  CONTINUE
      IF (NOUT.LE.3) GO TO 240
      P = 0.0
      DO 230 I=1,N1
        P = P + W(I)*ABS(X(I)-OUT(1,I))
230  CONTINUE
      IF (P.LE.E) GO TO 290
C CLOSED CURVE MAY BE EXPECTED
240  DXK2 = 1.0
      DO 250 I=1,N1
        DXK2 = DXK2 + BETA(I)**2
250  CONTINUE
      DXDT(K) = 1.0/SQRT(DXK2)*FLOAT(NDIR(K))
C DERIVATIVE OF INDEPENDENT VARIABLE X(K) WITH RESPECT TO ARC LENGTH
C OF SOLUTION LOCUS IS COMPUTED HERE
      H = HH
      DO 270 I=1,N1
        NDIR(I) = 1
        IF (I.EQ.K) GO TO 260
        DXDT(I) = BETA(I)*DXDT(K)
260  IF (DXDT(I).LT.0.0) NDIR(I) = -1
        IF (H*ABS(DXDT(I)).LE.HMAX(I)) GO TO 270
        MADMS = 0
        H = HMAX(I)/ABS(DXDT(I))
270  CONTINUE
      IF (NOUT.LE.KOUT+3) GO TO 280
      IF (H*ABS(DXDT(K)).LE.0.8*ABS(X(K)-OUT(1,K))) GO TO 280
      IF ((OUT(1,K)-X(K))*FLOAT(NDIR(K)).LE.0.0) GO TO 280
      MADMS = 0
      IF (H*ABS(DXDT(K)).LE.ABS(X(K)-OUT(1,K))) GO TO 280
      H = ABS(X(K)-OUT(1,K))/ABS(DXDT(K))
      KOUT = NOUT
280  CALL ADAMS(N, DXDT, MADMS, H, X, MXADMS)
      GO TO 70
290  WRITE (LW,99991)
      MAXOUT = -1
      RETURN
300  MAXOUT = -2
      RETURN
310  WRITE (LW,99990) (X(I),I=1,N1)
      MAXOUT = -3
      RETURN
99999 FORMAT (3X, 7H DERPAR, I3, 25H.INITIAL NEWTON ITERATION/22X,
* 11HX,ALFA,SQF=, 5F15.7/(33X, 5F15.7))
99998 FORMAT (/16H DERPAR OVERFLOW, I5, 2X, 18HINITIAL ITERATIONS/)
99997 FORMAT (3X, 39H DERPAR AFTER INITIAL NEWTON ITERATIONS/22X,
* 7HX,ALFA=, 4X, 5F15.7/(33X, 5F15.7))
99996 FORMAT (/6X, 45H DERPAR RESULTS (VARIABLE CHOSEN AS INDEPENDENDE,
* 18HNT IS MARKED BY *)/)
99995 FORMAT (/37H DERPAR NUMBER OF NEWTON CORRECTIONS=, I5,
* 31H IS NOT SUFFICIENT,ERROR OF X=, F15.7/)
99994 FORMAT (6X, 27H DERPAR RESULTS X,ALFA,SQF=)
99993 FORMAT (33X, 5(F15.7, A1))
99992 FORMAT (/28H DERPAR OUTPUT ARRAY IS FULL//)
99991 FORMAT (/36H DERPAR CLOSED CURVE MAY BE EXPECTED/)
99990 FORMAT (/48H DERPAR SINGULAR JACOBIAN MATRIX FOR X AND ALFA=,
* 5F12.6/(48X, 5F12.6))
      END
SUBROUTINE GAUSE(N, A, B, M, NN, MM, PREF, BETA, K)
      DIMENSION A(NN,MM), B(MM), PREF(MM), BETA(MM), Y(11), X(11),

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DER 1410
 DER 1420
 DER 1430
 DER 1440
 DER 1450
 DER 1460
 DER 1470
 DER 1480
 DER 1490
 DER 1500
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 DER 2230
 DER 2240
 DER 2250
 DER 2260
 DER 2270
 DER 2280
 DER 2290

GAU 10
 GAU 20

```

* IRR(11), IRK(11)
C SOLUTION OF N LINEAR EQUATIONS FOR N+1 UNKNOWNNS
C BASED ON GAUSSIAN ELIMINATION WITH PIVOTING.
C N - NUMBER OF EQUATIONS
C A - N X (N+1) MATRIX OF SYSTEM
C B - RIGHT-HAND SIDES
C M - IF (M.EQ.0) AFTER RETURN THEN RANK(A).LT.N
C PREF(I) - PREFERENCE NUMBER FOR X(I) TO BE INDEPENDENT VARIABLE,
C 0.0.LE.PREF(I).LE.1.0, THE LOWER IS PREF(I) THE HIGHER IS
C PREFERENCE OF X(I).
C BETA(I) - COEFFICIENTS IN EXPLICIT DEPENDENCES OBTAINED IN FORM -
C X(I)=B(I)+BETA(I)*X(K), I.NE.K.
C K - RESULTING INDEX OF INDEPENDENT VARIABLE
  N1 = N + 1
  ID = 1
  M = 1
  DO 10 I=1,N1
    IRR(I) = 0
    IRR(I) = 0
10 CONTINUE
20 IR = 1
  IS = 1
  AMAX = 0.0
  DO 60 I=1,N
    IF (IRR(I)) 60, 30, 60
30 DO 50 J=1,N1
  P = PREF(J)*ABS(A(I,J))
  IF (P-AMAX) 50, 50, 40
40 IR = I
  IS = J
  AMAX = P
50 CONTINUE
60 CONTINUE
  IF (AMAX.NE.0.0) GO TO 70
C THIS CONDITION FOR SINGULARITY OF MATRIX MUST BE SPECIFIED
C MORE EXACTLY WITH RESPECT TO COMPUTER ACTUALLY USED
  M = 0
  GO TO 150
70 IRR(IR) = IS
  DO 90 I=1,N
    IF (I.EQ.IR .OR. A(I,IS).EQ.0.0) GO TO 90
    P = A(I,IS)/A(IR,IS)
    DO 80 J=1,N1
      A(I,J) = A(I,J) - P*A(IR,J)
80 CONTINUE
    A(I,IS) = 0.0
    B(I) = B(I) - P*B(IR)
90 CONTINUE
  ID = ID + 1
  IF (ID.LE.N) GO TO 20
  DO 100 I=1,N
    IR = IRR(I)
    X(IR) = B(I)/A(I,IR)
    IRK(IR) = 1
100 CONTINUE
  DO 110 K=1,N1
    IF (IRK(K).EQ.0) GO TO 120
110 CONTINUE
120 DO 130 I=1,N
  IR = IRR(I)
  Y(IR) = -A(I,K)/A(I,IR)
130 CONTINUE
  DO 140 I=1,N1
    B(I) = X(I)
    BETA(I) = Y(I)
140 CONTINUE
  B(K) = 0.0
  BETA(K) = 0.0
150 RETURN
  END

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```

SUBROUTINE ADAMS(N, D, MADMS, H, X, MXADMS)
DIMENSION DER(4,11), X(11), D(11)
C ADAMS-BASHFORTH METHODS
  N1 = N + 1
  DO 20 I=1,3
    DO 10 J=1,N1
      DER(I+1,J) = DER(I,J)
10 CONTINUE
20 CONTINUE
  MADMS = MADMS + 1
  IF (MADMS.GT.MXADMS) MADMS = MXADMS
  IF (MADMS.GT.4) MADMS = 4
  DO 70 I=1,N1
    DER(1,I) = D(I)
    GO TO (30, 40, 50, 60), MADMS
30 X(I) = X(I) + H*DER(1,I)
  GO TO 70
40 X(I) = X(I) + 0.5*H*(3.0*DER(1,I)-DER(2,I))
  GO TO 70

```

GAU 30
GAU 40
GAU 50
GAU 60
GAU 70
GAU 80
GAU 90
GAU 100
GAU 110
GAU 120
GAU 130
GAU 140
GAU 150
GAU 160
GAU 170
GAU 180
GAU 190
GAU 200
GAU 210
GAU 220
GAU 230
GAU 240
GAU 250
GAU 260
GAU 270
GAU 280
GAU 290
GAU 300
GAU 310
GAU 320
GAU 330
GAU 340
GAU 350
GAU 360
GAU 370
GAU 380
GAU 390
GAU 400
GAU 410
GAU 420
GAU 430
GAU 440
GAU 450
GAU 460
GAU 470
GAU 480
GAU 490
GAU 500
GAU 510
GAU 520
GAU 530
GAU 540
GAU 550
GAU 560
GAU 570
GAU 580
GAU 590
GAU 600
GAU 610
GAU 620
GAU 630
GAU 640
GAU 650
GAU 660
GAU 670
GAU 680
GAU 690
GAU 700
GAU 710
GAU 720

ADA 10
ADA 20
ADA 30
ADA 40
ADA 50
ADA 60
ADA 70
ADA 80
ADA 90
ADA 100
ADA 110
ADA 120
ADA 130
ADA 140
ADA 150
ADA 160
ADA 170
ADA 180
ADA 190

COLLECTED ALGORITHMS (cont.)

50	X(I) = X(I) + H*(23.0*DER(1,I)-16.0*DER(2,I)+5.0*DER(3,I))/	ADA	200
	* 12.0	ADA	210
	GO TO 70	ADA	220
60	X(I) = X(I) + H*(55.0*DER(1,I)-59.0*DER(2,I)+37.0*DER(3,I)	ADA	230
	* -9.0*DER(4,I))/24.0	ADA	240
70	CONTINUE	ADA	250
	RETURN	ADA	260
	END	ADA	270

ALGORITHM 503

An Automatic Program for Fredholm Integral Equations of the Second Kind [D5]

KENDALL ATKINSON

University of Iowa

Key Words and Phrases: numerical analysis, linear integral equations, automatic algorithm, Nyström method

CR Categories: 5.18

Language: Fortran

DESCRIPTION

The two algorithms given here, *IESIMP* and *IEGAUS*, are a complement of [1] where the theoretical development is described.

REFERENCES

1. ATKINSON, K. An automatic program for Fredholm integral equations of the second kind. *ACM Trans. Math. Software* 2, 2 (June 1976), 154-171.

ALGORITHMS

```

      SUBROUTINE IESIMP(KERNEL, RHFCN, A, B, NZ, EPS, IFLAG, X,      IES  10
      * NUPPER, MUPPER, W, IFLGR2, IER)                          IES  20
C THE INTEGRAL EQUATION BEING SOLVED IS                          IES  30
C           B                                                    IES  40
C           X(S) - INT KERNEL(S,T)*X(T)*DT = RHFCN(S)          IES  50
C           A                                                    IES  60
C THE METHOD BEING USED IS BASED ON THE NYSTROM METHOD WITH      IES  70
C SIMPSONS RULE, WITH AN ITERATIVE TECHNIQUE OF SOLUTION FOR THE IES  80
C RESULTING LINEAR SYSTEM.                                       IES  90
C KERNEL      THESE ARE DOUBLE PRECISION FUNCTIONS OF TWO AND ONE IES 100
C RHFCN      VARIABLES, RESPECTIVELY. THEY MUST BE DECLARED IN AN IES 110
C           EXTERNAL STATEMENT IN THE PROGRAM CALLING IESIMP.   IES 120
C NZ         THE INITIAL VALUE OF N IN THE PROGRAM. N IS THE ORDER IES 130
C           OF AN INVERSE MATRIX WHICH IS BEING USED TO         IES 140
C           ITERATIVELY SOLVE A LARGER SYSTEM OF ORDER M, WHICH IES 150
C           APPROXIMATES THE ABOVE INTEGRAL EQUATION. THE CHOICE IES 160
C           OF NZ MUST ALWAYS BE ODD AND GREATER THAN TWO. FOR A IES 170
C           DEFAULT CHOICE, SET NZ=3.                             IES 180
C           ON COMPLETION OF THE PROGRAM, NZ IS SET EQUAL TO THE IES 190
C           FINAL NUMBER OF NODE POINTS USED IN OBTAINING THE   IES 200
C           SOLUTION. THE ARRAYS W AND X WILL CONTAIN THE NZ     IES 210
C           FINAL NODE POINTS AND CORRESPONDING SOLUTION        IES 220
C           VALUES.                                             IES 230
C EPS       THE DESIRED ERROR, INTERPRETED ACCORDING TO IFLAG.   IES 240

```

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ACM Transactions on Mathematical Software, Vol. 2, No. 2, June 1976, Pages 196-199.

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C          THIS VARIABLE IS CHANGED ON COMPLETION OF THE          IES 250
C          PROGRAM. SEE THE DISCUSSION OF IFLAG AND IER FOR MORE  IES 260
C          INFORMATION.                                           IES 270
C IFLAG =0 EPS IS TO BE INTERPRETED AS AN ABSOLUTE ERROR.        IES 280
C          =1 EPS IS TO BE INTERPRETED AS A RELATIVE ERROR.      IES 290
C X        THE ANSWER TO THE INTEGRAL EQUATION, EVALUATED AT THE  IES 300
C          FINAL SET OF NODE POINTS, WHICH ARE STORED IN W. THE  IES 310
C          DIMENSION OF X SHOULD BE AT LEAST MUPPER.            IES 320
C NUPPER   AN UPPER LIMIT ON THE VARIABLE N IN THIS PROGRAM. N   IES 330
C          IS THE ORDER OF APPROXIMATE INVERSE WHICH IS BEING   IES 340
C          USED TO ITERATIVELY SOLVE A LARGER SYSTEM OF ORDER M. IES 350
C MUPPER   AN UPPER LIMIT ON THE VARIABLE M IN THE PROGRAM.      IES 360
C          M IS THE NUMBER OF NODE POINTS BEING USED IN THE     IES 370
C          SIMPSON RULE APPROXIMATION OF THE INTEGRAL OPERATOR.  IES 380
C          N AND M ARE ALWAYS ODD INTEGERS, GREATER THAN TWO.    IES 390
C W        THIS IS TEMPORARY WORKING STORAGE. IT SHOULD HAVE     IES 400
C          DIMENSION 4*NU*NU+10*NU+9*MU, WITH NU=NUPPER,        IES 410
C          MU=MUPPER. ON EXIT, W WILL CONTAIN THE FINAL CHOICE   IES 420
C          OF NODE POINTS,CORRESPONDING TO THE APPROXIMATE      IES 430
C          SOLUTION VALUES STORED IN X.                         IES 440
C IFLGR2 =0 THE NORMAL SETTING.                                  IES 450
C          =1 IF THE INTEGRAND K(S,T)*X(T), REGARDED AS A FUNCTION IES 460
C          OF T, IS KNOWN TO BE PERIODIC ALONG WITH A NUMBER OF  IES 470
C          ITS DERIVATIVES ON THE INTERVAL (A,B), FOR ALL S IN   IES 480
C          THE INTERVAL, THEN THE ERROR ESTIMATES WILL BE MORE  IES 490
C          ACCURATE IF IFLGR2=1. THIS WILL GENERALLY RESULT IN  IES 500
C          MORE RAPID CONVERGENCE OF THE PROGRAM.                IES 510
C IER =0   THIS ERROR COMPLETION CODE MEANS THAT THE ERROR TEST  IES 520
C          WAS SATISFIED. THE PREDICTED ERROR IS STORED IN EPS.  IES 530
C          =1 THE ERROR TEST WAS NOT SATISFIED. THE PREDICTED ERROR IES 540
C          IS IN EPS.                                           IES 550
C          =2 THE ERROR TEST WAS NOT SATISFIED. THE VARIABLE EPS  IES 560
C          HAS BEEN SET TO ZERO.                                 IES 570
C          REGARDLESS OF THE VALUE OF IER, THE MOST ACCURATE     IES 580
C          SOLUTION OBTAINED IS STORED IN X.                     IES 590
C IESIMP IS PRESENTLY LIMITED TO NUPPER .LE. 100. THIS IS ENTIRELY IES 600
C DUE TO A RESTRICTION IN THE PROGRAM LINSYS. TO REMOVE THIS   IES 610
C RESTRICTION, CHANGE THE DIMENSION STATEMENT IN THE           IES 620
C COMMON/LINEAR/ STATEMENT IN LINSYS.                          IES 630
C          DOUBLE PRECISION KERNEL, RHFCN, A, B, EPS, X, W, CUTOFF, IES 640
C          * RATIO, ROOTRT                                       IES 650
C          DIMENSION X(1), W(1)                                  IES 660
C          EXTERNAL KERNEL, RHFCN                               IES 670
C *****                                                    IES 680
C          *                                                    IES 690
C          COMMON /INFO/ R1, R2, NFINAL                          IES 700
C          *                                                    IES 710
C          THE VARIABLES IN COMMON/INFO/ GIVE ADDITIONAL INFORMATION ABOUT* IES 720
C          THE FUNCTIONING OF IESIMP. R1 GIVES THE FINAL ITERATIVE RATE * IES 730
C          OF CONVERGENCE FOR SOLVING THE LINEAR SYSTEMS OF ORDER M. R2 * IES 740
C          GIVES THE RATE OF CONVERGENCE OF THE NYSTROM METHOD. FOR A * IES 750
C          SMOOTH KERNEL FUNCTION AND SMOOTH SOLUTION FUNCTION X(S), THE * IES 760
C          VALUE OF R2 SHOULD BE 0.0625 FOR SIMPSONS RULE. NFINAL GIVES * IES 770
C          THE FINAL VALUE OF N USED IN ITERATIVELY SOLVING THE LARGER * IES 780
C          SYSTEMS OF ORDER M. IF THE VALUE OF NFINAL EQUALS THE FINAL * IES 790
C          VALUE OF NZ, THEN ITERATION WAS NOT INVOKED SUCCESSFULLY. * IES 800
C          *                                                    IES 810
C *****                                                    IES 820
C          DATA CUTOFF /0.5D0/                                  IES 830
C *****                                                    IES 840
C          *                                                    IES 850
C          DATA RATIO /0.0625D0/, ROOTRT /0.25D0/             IES 860
C          *                                                    IES 870
C          THESE CONSTANTS DEPEND ON THE NUMERICAL INTEGRATION RULE * IES 880
C          BEING USED. SEE THE DISCUSSION OF IESIMP IN ORDER TO SET THEM * IES 890
C          PROPERLY WHEN CHANGING NUMERICAL INTEGRATION RULES. RATIO * IES 900
C          DEPENDS ON THE RATE OF CONVERGENCE OF THE RULE BEING USED. * IES 910
C          ROOTRT USUALLY EQUALS SQRT(RATIO).                   * IES 920
C          *                                                    IES 930
C *****                                                    IES 940
C          SET UP RELATIVE BASE ADDRESSES FOR THE VARIOUS ARRAYS INTO WHICH IES 950
C          W IS TO BE SPLIT.                                     IES 960
C          N = NUPPER                                           IES 970

```

```

M = MUPPER
NSQ = N*N
I1 = 1
I2 = I1 + NSQ
I3 = I2 + NSQ
I4 = I3 + (NSQ+N)/2
I5 = I4 + (NSQ+N)/2
I6 = I5 + M
I7 = I6 + M
I8 = I7 + N
I9 = I8 + N
I10 = I9 + M
I11 = I10 + N
I12 = I11 + M
I13 = I12 + M
I14 = I13 + M
I15 = I14 + N
I16 = I15 + M
I17 = I16 + M
I18 = I17 + N
I19 = I18 + M
I20 = I19 + 4*N
NHALF = (N+1)/2
CALL IESP(KERNEL, RHFCN, A, B, NZ, EPS, IFLAG, X, NUPPER,
* MUPPER, IFLGR2, IER, RATIO, ROOTRT, CUTOFF, NHALF, W, W(I1),
* W(I2), W(I3), W(I4), W(I5), W(I6), W(I7), W(I8), W(I9),
* W(I10), W(I11), W(I12), W(I13), W(I14), W(I15), W(I16),
* W(I17), W(I18), W(I19), W(I20))
RETURN
END
IES 980
IES 990
IES 1000
IES 1010
IES 1020
IES 1030
IES 1040
IES 1050
IES 1060
IES 1070
IES 1080
IES 1090
IES 1100
IES 1110
IES 1120
IES 1130
IES 1140
IES 1150
IES 1160
IES 1170
IES 1180
IES 1190
IES 1200
IES 1210
IES 1220
IES 1230
IES 1240
IES 1250
IES 1260
IES 1270

SUBROUTINE IESP(KERNEL, RHFCN, A, B, NZ, EPS, IFLAG, X, NUP,
* MUP, IFLGR2, IER, RATIO, ROOTRT, CUTOFF, NHALF, T, LUFAC,
* KMM, KMN, KNM, RHS, R, RH, DELN, TM, TN, XM, XMZ, WM, WN,
* OLDX, SAVE, XN, SAVE2, ASIDE, ASIDE3)
C THIS ROUTINE CONTROLS THE SOLUTION OF THE INTEGRAL EQUATION.
DOUBLE PRECISION KERNEL, RHFCN, A, B, EPS, X, RATIO, ROOTRT,
* CUTOFF, T, LUFAC, KMM, KMN, KNM, RHS, R, RH, DELN, TM, TN,
* XM, XMZ, WM, WN, OLDX, SAVE, SAVE2, XN, ASIDE, ASIDE2,
* ASIDE3, R2, R1RAT, NUMR2, DENR2, NORM, DMIN1, DSQRT, ERROR,
* TEMP, TEST, R1, DENR1, NUMR1, RT, DET
INTEGER OLDX, TWICE, FLAG
DIMENSION X(MUP), T(MUP), LUFAC(NUP,NUP), KMM(NUP,NUP),
* KMN(NUP,NHALF), KNM(NHALF,NUP), RHS(MUP), R(MUP), RH(NUP),
* DELN(NUP), TM(MUP), TN(NUP), XM(MUP), XMZ(MUP), WM(MUP),
* WN(NUP), OLDX(MUP), SAVE(MUP), XN(NUP), SAVE2(MUP),
* ASIDE(NUP,4), ASIDE2(5), ASIDE3(NUP,NUP)
COMMON /INFO/ R1, R2, N
EXTERNAL KERNEL, RHFCN
C *****
C
TWICE(KK) = 2*KK - 1
C
C
C THIS FUNCTION GIVES THE FORMULA FOR INCREASING THE NUMBER OF
C NODE POINTS. WITH ANOTHER INTEGRATION RULE, IT MAY BE
C NECESSARY TO CHANGE THE DEFINITION OF TWICE(KK).
C
C *****
C
INITIALIZATION.
IER = 0
LOOP = 1
N = NZ
R2 = 0.5D0
M = TWICE(N)
R1RAT = ROOTRT
C STAGE A. DIRECT SOLUTION OF LINEAR SYSTEM (I-KN)*XN=RHS, WHILE
C TRYING TO FIND A GOOD APPROXIMATE INVERSE TO IMPLEMENT THE
C ITERATIVE METHOD OF SOLUTION.
C CREATE NODES TN(I) AND WEIGHTS WN(I), I=1,...,N.
CALL WANDT(WN, TN, N, A, B)
C CREATE SYSTEM (I-KN)*XN=RHFCN.
IES 10
IES 20
IES 30
IES 40
IES 50
IES 60
IES 70
IES 80
IES 90
IES 100
IES 110
IES 120
IES 130
IES 140
IES 150
IES 160
IES 170
IES 180
IES 190
IES 200
IES 210
IES 220
IES 230
IES 240
IES 250
IES 260
IES 270
IES 280
IES 290
IES 300
IES 310
IES 320
IES 330
IES 340
IES 350
IES 360
IES 370
IES 380
IES 390
IES 400

```

```

DO 20 I=1,N
DO 10 J=1,N
LUFAC(T,I,J) = -WN(J)*KERNEL(TN(I),TN(J))
10 CONTINUE
XN(I) = RHFCN(TN(I))
LUFAC(T,I,I) = LUFAC(T,I,I) + 1.0D0
20 CONTINUE
GO TO 60
C THIS IS ENTRANCE FOR AN INCREASED VALUE OF N, USING PREVIOUSLY
C STORED VALUES IN KMM AND RHS.
30 DO 50 J=1,N
DO 40 I=1,N
LUFAC(T,I,J) = -KMM(I,J)
40 CONTINUE
WN(J) = WM(J)
TN(J) = TM(J)
XN(J) = RHS(J)
LUFAC(T,J,J) = LUFAC(T,J,J) + 1.0D0
50 CONTINUE
C THIS IS THE ENTRANCE WHEN ITERATION IN STAGE B FAILS AND WE NEED
C TO INCREASE N TO OBTAIN A BETTER ITERATIVE RATE.
60 CONTINUE
C SOLVE (I-KN)*XN=RHFCN AT ALL TN(I).ALSO OBTAIN THE LU
C DECOMPOSITION FOR LATER USE IN THE STAGE B ITERATIVE METHOD.
C *****
C CALL LINSYS(LUFAC, LUFAC, N, XN, XN, 1, DET, NUP)
C
C THIS LINEAR EQUATION SOLVER MAY BE REPLACED BY ANOTHER PROGRAM,*
C BUT CARE MUST BE TAKEN THAT THE NEW PROGRAM DOES THE SAME JOB *
C AS LINSYS. THE PROGRAM LINSYS IS ALSO REFERENCED IN THE *
C SUBROUTINE ITERT. *
C *****
IF (LOOP.EQ.1) GO TO 110
IF (LOOP.EQ.2) GO TO 80
C COMPUTE THE RATE OF CONVERGENCE, AND COMPARE WITH THE
C THEORETICAL RATIO.
NUMR2 = NORM(XN,OLDX,N,1)
R2 = DMIN1(0.5D0,NUMR2/DENR2)
IF ((R2.LT.RATIO) .AND. (IFLGR2.EQ.0)) R2 = RATIO
R1RAT = DMIN1(DSQRT(R2),ROOTRT)
IF (IFLGR2.EQ.0) R1RAT = ROOTRT
C CHECK FOR ERROR IN XN USING TEST INVOLVING R2 AND OLDX,ACCORDING
C TO THEORY FOR ASYMPTOTIC ERROR BOUNDS.
70 ERROR = (R2/(1.0D0-R2))*NUMR2
IF (IFLAG.EQ.1) ERROR = ERROR/NORM(XN,XN,N,0)
IF ((IFLGR2.EQ.1) .AND. (R2.LT.RATIO)) ERROR = 2.0*ERROR
IF (ERROR.LE.EPS) GO TO 90
DENR2 = NUMR2
GO TO 110
C ENTRANCE FOR LOOP=2.
80 NUMR2 = NORM(XN,OLDX,N,1)
DENR2 = 0.0D0
GO TO 70
C SET UP T,X,EPS,NZ FOR SUCCESSFUL RETURN.
90 DO 100 I=1,N
X(I) = XN(I)
T(I) = TN(I)
100 CONTINUE
EPS = ERROR
NZ = N
RETURN
C INITIALIZE FOR SOLVING (I-KM)*XM=RHFCN ITERATIVELY.
110 CALL WANDT(WM, TM, M, A, B)
FLAG = 0
CALL INTERP(TM, WM, XMZ, M, TN, WN, XN, N, KERNEL, RHFCN,
* RHS, KMM, NHALF, NUP)
DO 120 I=1,M
OLDX(I) = XMZ(I)
120 CONTINUE
CALL ITERT(KERNEL, RHFCN, N, TN, WN, M, TM, WM, XM, XMZ, KMM,
* KMM, KMM, RHS, LUFAC, R, RH, DELN, NUP, NHALF, FLAG)
DENR1 = NORM(XM,XMZ,M,1)
FLAG = 1
DO 130 I=1,M

```

```

IES 410
IES 420
IES 430
IES 440
IES 450
IES 460
IES 470
IES 480
IES 490
IES 500
IES 510
IES 520
IES 530
IES 540
IES 550
IES 560
IES 570
IES 580
IES 590
IES 600
IES 610
IES 620
IES 630
IES 640
IES 650
IES 660
IES 670
IES 680
IES 690
IES 700
IES 710
IES 720
IES 730
IES 740
IES 750
IES 760
IES 770
IES 780
IES 790
IES 800
IES 810
IES 820
IES 830
IES 840
IES 850
IES 860
IES 870
IES 880
IES 890
IES 900
IES 910
IES 920
IES 930
IES 940
IES 950
IES 960
IES 970
IES 980
IES 990
IES 1000
IES 1010
IES 1020
IES 1030
IES 1040
IES 1050
IES 1060
IES 1070
IES 1080
IES 1090
IES 1100
IES 1110
IES 1120
IES 1130
IES 1140
IES 1150
IES 1160

```

```

      XMZ(I) = XM(I)
130 CONTINUE
      CALL ITERT(KERNEL, RHFCN, N, TN, WN, M, TM, WM, XM, XMZ, KMM,
      * KMN, KNM, RHS, LUFAC, R, RH, DELN, NUP, NHALF, FLAG)
      NUMR1 = NORM(XM, XMZ, M, 1)
C CHECK ON THE SPEED OF CONVERGENCE OF ITERATIVE METHOD. IF IT IS
C SUFFICIENTLY RAPID, THEN FIX N AND GO TO STAGE B.
      R1 = NUMR1/DENR1
      IF (M.GT.NUP) GO TO 190
      IF (R1.LE.R1RAT) GO TO 150
C REINITIALIZE FOR SOLVING (I-KN)*XN=RHFCN AGAIN WITH A LARGER N.
140 N = M
      LOOP = LOOP + 1
      M = TWICE(N)
      GO TO 30
C SAVE INFORMATION IN CASE STAGE B ABORTS AT A LARGER VALUE OF M
C AND STAGE A HAS TO BE RETURNED TO.
150 DO 160 I=1,M
      ASIDE(I,1) = OLDX(I)
      ASIDE(I,2) = WM(I)
      ASIDE(I,3) = TM(I)
      ASIDE(I,4) = RHS(I)
160 CONTINUE
      ASIDE2(1) = LOOP
      ASIDE2(2) = M
      ASIDE2(3) = R2
      ASIDE2(4) = DENR2
      ASIDE2(5) = R1RAT
      DO 180 I=1,M
        DO 170 J=1,M
          ASIDE3(I,J) = KMM(I,J)
170 CONTINUE
180 CONTINUE
      GO TO 240
C STAGE B. ITERATIVE METHOD OF SOLUTION OF (I-KM)*XM=RHS.
190 IF (R1.LE.CUTOFF) GO TO 240
C IF ITERATES ARE CONVERGING VERY SLOWLY OR NOT AT ALL, THEN
C RETURN WITHOUT FURTHER ATTEMPTS TO LESSEN ERROR.
200 IF (LOOP.NE.1) GO TO 230
      EPS = 0.0D0
      IER = 2
210 DO 220 I=1,N
      X(I) = XN(I)
      T(I) = TN(I)
220 CONTINUE
      NZ = N
      RETURN
230 EPS = ERROR
      IER = 1
      GO TO 210
C TEST TO SEE IF THE CURRENT ITERATE XM IS SUFFICIENTLY ACCURATE
C COMPARED TO THE TRUE XM.
240 TEMP = NORM(XM, OLDX, M, 1)
      RT = DMIN1(RATIO, R2)
      TEST = ((1.0D0-R1)/R1)*(RT/(1.0D0-RT))*TEMP
      IF (NUMR1.LE.TEST) GO TO 320
C ITERATE NOT ACCURATE. INITIALIZE FOR COMPUTATION OF ANOTHER
C ITERATE.
      DENR1 = NUMR1
      DO 250 I=1,M
        XMZ(I) = XM(I)
250 CONTINUE
      CALL ITERT(KERNEL, RHFCN, N, TN, WN, M, TM, WM, XM, XMZ, KMM,
      * KMN, KNM, RHS, LUFAC, R, RH, DELN, NUP, NHALF, FLAG)
      NUMR1 = NORM(XM, XMZ, M, 1)
      R1 = NUMR1/DENR1
      IF (R1.LE.CUTOFF) GO TO 240
C THIS IS ENTRANCE FOR CASE WHERE ITERATION FAILS IN STAGE B.
C PARAMETERS MUST BE RESET FOR A RETURN TO STAGE A, OR IF N CANNOT
C BE INCREASED, FOR AN ERROR EXIT FROM IESIMP.
260 MNEW = ASIDE2(2)
      IF (MNEW.GT.NUP) GO TO 290
      IF (M.EQ.TWICE(N)) GO TO 140
      N = MNEW
      DO 280 J=1,N
        DO 270 I=1,N

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```

          LUFAC(T,I,J) = -ASIDE3(I,J)
270 CONTINUE
          OLDX(J) = ASIDE(J,1)
          WN(J) = ASIDE(J,2)
          TN(J) = ASIDE(J,3)
          XN(J) = ASIDE(J,4)
          LUFAC(J,J) = LUFAC(J,J) + 1.0D0
280 CONTINUE
          M = TWICE(N)
          LOOP = ASIDE2(1) + 1.0D0
          R2 = ASIDE2(3)
          DENR2 = ASIDE2(4)
          R1RAT = ASIDE2(5)
          GO TO 60
C ABORTIVE EXIT FROM STAGE B. N CANNOT BE INCREASED FURTHER, AND
C R1 IS NOT SUFFICIENTLY SMALL.
290 IF (M.EQ.TWICE(N)) GO TO 200
          DO 300 I=1,OLDM
              T(I) = SAVE(I)
300 CONTINUE
          NZ = OLDM
          IF (LOOP.EQ.1) GO TO 310
          EPS = ERROR
          IER = 1
          RETURN
310 EPS = 0.0
          IER = 2
          RETURN
C AN ACCURATE VALUE OF XM HAS BEEN OBTAINED. R2 IS TO BE COMPUTED
C AND COMPARED TO RATIO. THEN THE ERROR IN XM IS TO BE ESTIMATED.
320 IF (LOOP.EQ.1) GO TO 340
          NUMR2 = TEMP
          R2 = DMIN1(NUMR2/DENR2,0.5D0)
          IF ((R2.LT.RATIO) .AND. (IFLGR2.EQ.0)) R2 = RATIO
          DENR2 = NUMR2
330 ERROR = (R2/(1.0D0-R2))*TEMP
          IF (IFLAG.EQ.1) ERROR = ERROR/NORM(XM,XM,M,0)
          IF ((IFLGR2.EQ.1) .AND. (R2.LT.RATIO)) ERROR = 2.0*ERROR
          IF (ERROR.LE.EPS) GO TO 400
          MNEW = TWICE(M)
          IF (MNEW.LE.MUP) GO TO 350
          IER = 1
          GO TO 400
340 DENR2 = TEMP
          LOOP = 2
          GO TO 330
C ERROR NOT SUFFICIENTLY SMALL. M IS INCREASED AND TWO MORE
C ITERATES ARE COMPUTED WITH A NEW M.
350 DO 360 I=1,M
          X(I) = XM(I)
360 CONTINUE
          OLDM = M
          M = MNEW
          DO 370 I=1,OLDM
              SAVE2(I) = WM(I)
              SAVE(I) = TM(I)
370 CONTINUE
          CALL WANDT(WM, TM, M, A, B)
          FLAG = 0
          CALL INTERP(TM, WM, XMZ, M, SAVE, SAVE2, XM, OLDM, KERNEL,
* RHFCN, RHS, KMN, NHALF, NUP)
          DO 380 I=1,M
              OLDX(I) = XMZ(I)
380 CONTINUE
          CALL ITERT(KERNEL, RHFCN, N, TN, WN, M, TM, WM, XM, XMZ, KMM,
* KMN, KNM, RHS, LUFAC, R, RH, DELN, NUP, NHALF, FLAG)
          DENR1 = NORM(XM,XMZ,M,1)
          FLAG = 1
          DO 390 I=1,M
              XMZ(I) = XM(I)
390 CONTINUE
          CALL ITERT(KERNEL, RHFCN, N, TN, WN, M, TM, WM, XM, XMZ, KMM,
* KMN, KNM, RHS, LUFAC, R, RH, DELN, NUP, NHALF, FLAG)
          NUMR1 = NORM(XM,XMZ,M,1)
          R1 = NUMR1/DENR1
          IF (R1.LE.CUTOFF) GO TO 240

```

IES 1930
 IES 1940
 IES 1950
 IES 1960
 IES 1970
 IES 1980
 IES 1990
 IES 2000
 IES 2010
 IES 2020
 IES 2030
 IES 2040
 IES 2050
 IES 2060
 IES 2070
 IES 2080
 IES 2090
 IES 2100
 IES 2110
 IES 2120
 IES 2130
 IES 2140
 IES 2150
 IES 2160
 IES 2170
 IES 2180
 IES 2190
 IES 2200
 IES 2210
 IES 2220
 IES 2230
 IES 2240
 IES 2250
 IES 2260
 IES 2270
 IES 2280
 IES 2290
 IES 2300
 IES 2310
 IES 2320
 IES 2330
 IES 2340
 IES 2350
 IES 2360
 IES 2370
 IES 2380
 IES 2390
 IES 2400
 IES 2410
 IES 2420
 IES 2430
 IES 2440
 IES 2450
 IES 2460
 IES 2470
 IES 2480
 IES 2490
 IES 2500
 IES 2510
 IES 2520
 IES 2530
 IES 2540
 IES 2550
 IES 2560
 IES 2570
 IES 2580
 IES 2590
 IES 2600
 IES 2610
 IES 2620
 IES 2630
 IES 2640
 IES 2650
 IES 2660
 IES 2670
 IES 2680

GO TO 260	IES 2690
400 DO 410 I=1,M	IES 2700
X(I) = XM(I)	IES 2710
T(I) = TM(I)	IES 2720
410 CONTINUE	IES 2730
NZ = M	IES 2740
EPS = ERROR	IES 2750
RETURN	IES 2760
END	IES 2770
SUBROUTINE WANDT(W, T, N, A, B)	WAN 10
C THIS ROUTINE CALCULATES THE WEIGHTS AND NODE POINTS FOR	WAN 20
C SIMPSONS RULE ON (A,B), WITH N EVENLY SPACED NODES.	WAN 30
DOUBLE PRECISION W, T, A, B, FN, FI, TEMP, H	WAN 40
DIMENSION W(N), T(N)	WAN 50
C CALCULATE NODE POINTS.	WAN 60
FN = N	WAN 70
H = (B-A)/(FN-1.0D0)	WAN 80
DO 10 I=1,N	WAN 90
FI = I	WAN 100
T(I) = A + (FI-1.0D0)*H	WAN 110
10 CONTINUE	WAN 120
C CALCULATE WEIGHTS.	WAN 130
W(1) = H/3.0D0	WAN 140
W(N) = H/3.0D0	WAN 150
NM1 = N - 1	WAN 160
TEMP = 4.0D0*H/3.0D0	WAN 170
DO 20 I=2,NM1,2	WAN 180
W(I) = TEMP	WAN 190
20 CONTINUE	WAN 200
IF (N.EQ.3) RETURN	WAN 210
TEMP = 2.0D0*H/3.0D0	WAN 220
NM2 = N - 2	WAN 230
DO 30 I=3,NM2,2	WAN 240
W(I) = TEMP	WAN 250
30 CONTINUE	WAN 260
RETURN	WAN 270
END	WAN 280
SUBROUTINE INTERP(TM, WM, XM, M, TN, WN, XN, N, KERNEL,	INT 10
* RHFCN, RHS, KMN, NHALF, NUP)	INT 20
C USE THE VALUES OF XN(I), I=1,...,N, TO CALCULATE THE NYSTROM	INT 30
C INTERPOLATES XM(I), I=1,...,M.	INT 40
DOUBLE PRECISION KERNEL, RHFCN, TM, WM, XM, TN, WN, XN, RHS,	INT 50
* KMN	INT 60
DIMENSION TM(M), WM(M), XM(M), TN(N), WN(N), XN(N), RHS(M),	INT 70
* KMN(NUP,NHALF)	INT 80
IF (M.GT.NUP) GO TO 60	INT 90
C SINCE M .LE. NUPPER, SAVE K(TM(I),TN(J))=KMN(I,J) AND	INT 100
C RHS(I)=RHFCN(TM(I)) FOR LATER USE IN ITERT.	INT 110
DO 20 I=1,M	INT 120
DO 10 J=1,N	INT 130
KMN(I,J) = WN(J)*KERNEL(TM(I),TN(J))	INT 140
10 CONTINUE	INT 150
20 CONTINUE	INT 160
DO 30 I=1,M	INT 170
RHS(I) = RHFCN(TM(I))	INT 180
XM(I) = RHS(I)	INT 190
30 CONTINUE	INT 200
C CALCULATE NYSTROM INTERPOLATING FORMULA.	INT 210
DO 50 I=1,M	INT 220
DO 40 J=1,N	INT 230
XM(I) = XM(I) + KMN(I,J)*XN(J)	INT 240
40 CONTINUE	INT 250
50 CONTINUE	INT 260
RETURN	INT 270
C M .GT. NUPPER, SO SAVE JUST RHS(I) FOR LATER USE IN ITERT.	INT 280
C CALCULATE NYSTROM INTERPOLATING FORMULA.	INT 290
60 DO 80 I=1,M	INT 300
RHS(I) = RHFCN(TM(I))	INT 310
XM(I) = RHS(I)	INT 320
DO 70 J=1,N	INT 330

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      XM(I) = XM(I) + WN(J)*KERNEL(TM(I),TN(J))*XN(J)      INT 340
70  CONTINUE                                             INT 350
80  CONTINUE                                             INT 360
      RETURN                                             INT 370
      END                                               INT 380

      SUBROUTINE ITERT(KERNEL, RHFCN, N, TN, WN, M, TM, WM, XM,
*  XMZ, KMM, KMN, KNM, RHS, LUFAC, R, RH, DELN, NUP, NHALF,
*  IFLG)
C THIS ROUTINE CALCULATES ONE ITERATE XM GIVEN THE INITIAL GUESS
C XMZ. THE ROUTINE IS DIVIDED ACCORDING TO WHETHER OR NOT
C M .GT. NUPPER.
      DOUBLE PRECISION KERNEL, RHFCN, TN, WN, TM, WM, XM, XMZ, KMM,
*  KMN, KNM, RHS, LUFAC, R, RH, DELN, SUM, DET
      DIMENSION TN(N), WN(N), TM(M), WM(M), XM(M), XMZ(M),
*  KMM(NUP,NUP), KMN(NUP,NHALF), KNM(NHALF,NUP), RHS(M),
*  LUFAC(NUP,NUP), R(M), RH(N), DELN(N)
C M .GT. NUPPER MEANS THAT THE MATRICES KMM,KMN,KNM CAN NO LONGER
C BE STORED DUE TO LACK OF SPACE.
      IF (M.GT.NUP) GO TO 120
      IF (IFLG.EQ.1) GO TO 40
C IF IFLG=0, THEN THE MATRICES KMM AND KNM MUST BE COMPUTED
C AND STORED.
      DO 30 J=1,M
        DO 10 I=1,M
          KMM(I,J) = WM(J)*KERNEL(TM(I),TM(J))
10    CONTINUE
        DO 20 I=1,N
          KNM(I,J) = WM(J)*KERNEL(TN(I),TM(J))
20    CONTINUE
30    CONTINUE
C COMPUTE RESIDUALS R(I)=RHFCN(TM(I))-XMZ(I)+KM(TM(I))*XMZ(I)
40  DO 60 I=1,M
      SUM = 0.0D0
      DO 50 J=1,M
        SUM = SUM + KMM(I,J)*XMZ(J)
50    CONTINUE
      R(I) = RHS(I) - (XMZ(I)-SUM)
60  CONTINUE
C COMPUTE RH=KM*R AT ALL TN(I).
      DO 80 I=1,N
        RH(I) = 0.0D0
        DO 70 J=1,M
          RH(I) = RH(I) + KNM(I,J)*R(J)
70    CONTINUE
80    CONTINUE
C CALCULATE DELN=((I-KN)**(-1))*KM*R AT ALL TN(I).
C *****
C CALL LINSYS(LUFAC, LUFAC, N, RH, DELN, 3, DET, NUP)
C
C SEE THE ORIGINAL REFERENCE IN IESP.
C *****
C CALCULATE NEW XM.
      DO 110 I=1,M
        SUM = 0.0D0
        DO 90 J=1,M
          SUM = SUM + KMM(I,J)*R(J)
90    CONTINUE
        DO 100 J=1,N
          SUM = SUM + KMN(I,J)*DELN(J)
100   CONTINUE
        XM(I) = SUM + R(I) + XMZ(I)
110  CONTINUE
      RETURN
C ENTRANCE WHEN M .GT. NUP.
C CALCULATE RESIDUALS.
120 DO 140 I=1,M
      SUM = 0.0D0
      DO 130 J=1,M
        SUM = SUM + WM(J)*KERNEL(TM(I),TM(J))*XMZ(J)
130  CONTINUE
      R(I) = RHS(I) - (XMZ(I)-SUM)

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140	CONTINUE	ITE	690
C	CALCULATE RH=KM*R.	ITE	700
	DO 160 I=1,N	ITE	710
	RH(I) = 0.0D0	ITE	720
	DO 150 J=1,M	ITE	730
	RH(I) = RH(I) + WM(J)*KERNEL(TN(I),TM(J))*R(J)	ITE	740
	150 CONTINUE	ITE	750
	160 CONTINUE	ITE	760
C	*****	ITE	770
C		*	ITE 780
	CALL LINSYS(LUFACT, LUFACT, N, RH, DELN, 3, DET, NUP)	ITE	790
C		*	ITE 800
C	SEE THE ORIGINAL REFERENCE IN IESP.	*	ITE 810
C		*	ITE 820
C	*****	ITE	830
C	CALCULATE XM.	ITE	840
	DO 190 I=1,M	ITE	850
	SUM = 0.0D0	ITE	860
	DO 170 J=1,M	ITE	870
	SUM = SUM + WM(J)*KERNEL(TM(I),TM(J))*R(J)	ITE	880
170	CONTINUE	ITE	890
	DO 180 J=1,N	ITE	900
	SUM = SUM + WN(J)*KERNEL(TM(I),TN(J))*DELN(J)	ITE	910
180	CONTINUE	ITE	920
	XM(I) = SUM + R(I) + XMZ(I)	ITE	930
190	CONTINUE	ITE	940
	RETURN	ITE	950
	END	ITE	960
	DOUBLE PRECISION FUNCTION NORM(X, Y, N, IFLAG)	NOR	10
C	IFLAG=0 CALCULATE THE MAXIMUM NORM OF X.	NOR	20
C	IFLAG=1 CALCULATE THE MAXIMUM NORM OF X-Y.	NOR	30
	DOUBLE PRECISION X, Y, DMAX1, DABS	NOR	40
	DIMENSION X(N), Y(N)	NOR	50
	IF (IFLAG.EQ.1) GO TO 20	NOR	60
C	FIND THE NORM OF X.	NOR	70
	NORM = 0.0D0	NOR	80
	DO 10 I=1,N	NOR	90
	NORM = DMAX1(NORM,DABS(X(I)))	NOR	100
10	CONTINUE	NOR	110
	RETURN	NOR	120
C	FIND THE NORM OF X-Y.	NOR	130
	20 NORM = 0.0D0	NOR	140
	DO 30 I=1,N	NOR	150
	NORM = DMAX1(NORM,DABS(X(I)-Y(I)))	NOR	160
30	CONTINUE	NOR	170
	RETURN	NOR	180
	END	NOR	190
	SUBROUTINE LINSYS(A, D, N, B, X, OPTION, DET, MACHIN)	LIN	10
C	SOLVE A*X=B, ORDER(A)=N, DIMENSION OF A=MACHIN.	LIN	20
C	OPTION=1 SOLVE A*X=B, LEAVE THE LU DECOMPOSITION A=L*U IN D	LIN	30
C	AND THE PIVOTS IN PIVOT. THE ANSWERS ARE LEFT IN X.	LIN	40
C	IT IS PERMISSABLE TO LET B=X AND D=A, BUT THEN THE	LIN	50
C	ORIGINAL CONTENTS OF A AND B ARE LOST.	LIN	60
C	OPTION=2 CALCULATE DECOMPOSITION A=L*U INCLUDING PIVOTS. SOLVE	LIN	70
C	A*X=B, AND THEN CALCULATE THE RESIDUAL AND ONE	LIN	80
C	CORRECTION. THE CORRECTIONS ARE LEFT IN R, THE NEW	LIN	90
C	VALUE OF X1 IN X, THE RELATIVE ERROR	LIN	100
C	NORM(X0-X1)/NORM(X1)	LIN	110
C	IN THE VARIABLE ERROR, AND THE RELATIVE RESIDUAL	LIN	120
C	NORM(RESIDUAL)/NORM(B)	LIN	130
C	IN THE VARIABLE RELRSD. THESE VALUES CAN BE OBTAINED	LIN	140
C	USING THE COMMON/LINEAR/ GIVEN BELOW.	LIN	150
C	OPTION=3 SAME AS OPTION=1, EXCEPT A=L*U IS KNOWN AND STORED	LIN	160
C	IN D.	LIN	170
C	OPTION=4 SAME AS OPTION=2, EXCEPT A=L*U IS KNOWN AND STORED	LIN	180
C	IN D.	LIN	190
C	THE DECOMPOSITION OF A INTO L*U USES SCALED PARTIAL PIVOTING IN	LIN	200
C	THE COLUMNS. FOR OPTIONS 1 AND 2, THE DETERMINANT OF A IS	LIN	210
C	CALCULATED AND STORED IN DET. IF DET=0, THEN THE ANSWERS ARE	LIN	220
C	NONSENSE, D AND X DO NOT CONTAIN USEFUL INFORMATION, AND AND A	LIN	230

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C AND B ARE LEFT UNDISTURBED (UNLESS D=A OR X=B).          LIN 240
C LINSYS IS PRESENTLY LIMITED TO N .LE. 100. TO REMOVE THIS LIN 250
C RESTRICTION,CHANGE THE DIMENSION STATEMENT FOR THE ARRAYS R, LIN 260
C SCALE, AND PIVOT, WHICH ARE GIVEN IN COMMON/LINEAR/.    LIN 270
  INTEGER OPTION, PIVOT                                  LIN 280
  DOUBLE PRECISION NORMX, NORME, NORMB, NORMR, A, D, B, X, R, LIN 290
  * SCALE, ERROR, RELRSD, DET, C, TEMP, DMAX1, DABS, SUM   LIN 300
  DIMENSION A(MACHIN,MACHIN), D(MACHIN,MACHIN), B(N), X(N) LIN 310
  COMMON /LINEAR/ R(100), SCALE(100), ERROR, RELRSD, PIVOT(100) LIN 320
  ISWIT = 1                                              LIN 330
  IF (OPTION.GT.2) GO TO 110                             LIN 340
C PRODUCE LU DECOMPOSITION OF A AND DET(A)                LIN 350
  DET = 0.0D0                                           LIN 360
  DO 20 I=1,N                                           LIN 370
    DO 10 J=1,N                                          LIN 380
      D(I,J) = A(I,J)                                  LIN 390
    10 CONTINUE                                         LIN 400
  20 CONTINUE                                           LIN 410
C PRODUCE SCALING FACTORS FOR SCALED PARTIAL PIVOTING.   LIN 420
  DO 40 I=1,N                                           LIN 430
    SCALE(I) = 0.0D0                                     LIN 440
    DO 30 J=1,N                                          LIN 450
      SCALE(I) = DMAX1(SCALE(I),DABS(D(I,J)))           LIN 460
    30 CONTINUE                                         LIN 470
    IF (SCALE(I).EQ.0.0D0) RETURN                       LIN 480
  40 CONTINUE                                           LIN 490
  DET = 1.0D0                                           LIN 500
  NM1 = N - 1                                           LIN 510
C BEGIN GAUSSIAN ELIMINATION.                             LIN 520
  DO 100 K=1,NM1                                        LIN 530
C SELECT PIVOT ROW.                                       LIN 540
  C = DABS(D(K,K))/SCALE(K)                             LIN 550
  INDEX = K                                             LIN 560
  KP1 = K + 1                                          LIN 570
  DO 50 I=KP1,N                                         LIN 580
    TEMP = DABS(D(I,K))/SCALE(I)                       LIN 590
    IF (TEMP.LE.C) GO TO 50                             LIN 600
    INDEX = I                                           LIN 610
    C = TEMP                                            LIN 620
  50 CONTINUE                                           LIN 630
  PIVOT(K) = INDEX                                     LIN 640
  IF (INDEX.EQ.K) GO TO 70                              LIN 650
C SWITCH ROWS OF D.                                       LIN 660
  DO 60 J=K,N                                           LIN 670
    TEMP = D(K,J)                                       LIN 680
    D(K,J) = D(INDEX,J)                                 LIN 690
    D(INDEX,J) = TEMP                                   LIN 700
  60 CONTINUE                                           LIN 710
  TEMP = SCALE(K)                                       LIN 720
  SCALE(K) = SCALE(INDEX)                               LIN 730
  SCALE(INDEX) = TEMP                                   LIN 740
  DET = -DET                                           LIN 750
  70 DET = DET*D(K,K)                                   LIN 760
C ELIMINATE UNKNOWN =K FROM BELOW DIAGONAL IN COLUMN K. LIN 770
  IF (DET.EQ.0.0D0) RETURN                              LIN 780
  DO 90 I=KP1,N                                         LIN 790
    D(I,K) = D(I,K)/D(K,K)                              LIN 800
    TEMP = D(I,K)                                       LIN 810
    DO 80 J=KP1,N                                       LIN 820
      D(I,J) = D(I,J) - TEMP*D(K,J)                   LIN 830
    80 CONTINUE                                         LIN 840
  90 CONTINUE                                           LIN 850
100 CONTINUE                                           LIN 860
  DET = DET*D(N,N)                                       LIN 870
  IF (DET.EQ.0.0D0) RETURN                              LIN 880
C THE DECOMPOSITION A=P*L*U IS COMPLETE.                 LIN 890
C BEGIN SOLUTION OF LINEAR SYSTEM P*L*U*X=B             LIN 900
C SET R AND X FOR SOLUTION OF A*X=B.                   LIN 910
  110 DO 120 I=1,N                                       LIN 920
    R(I) = B(I)                                         LIN 930
    X(I) = 0.0D0                                        LIN 940
  120 CONTINUE                                           LIN 950
  GO TO 170                                             LIN 960
C SET R=B-A*X FOR COMPUTATION OF CORRECTION.           LIN 970
  130 DO 150 I=1,N                                       LIN 980
    SUM = 0.0D0                                         LIN 990

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      DO 140 J=1,N
          SUM = SUM + A(I,J)*X(J)
140  CONTINUE
      R(I) = B(I) - SUM
150  CONTINUE
      NORMB = 0.0D0
      NORMR = 0.0D0
      DO 160 I=1,N
          NORMB = DMAX1(NORMB,DABS(B(I)))
          NORMR = DMAX1(NORMR,DABS(R(I)))
160  CONTINUE
      RELRSD = NORMR/NORMB
      ISWIT = 2
C SOLVE P*L*Z=R AND STORE IN R.
170  NM1 = N - 1
      DO 200 K=1,NM1
          IF (PIVOT(K).EQ.K) GO TO 180
          KPIV = PIVOT(K)
          TEMP = R(K)
          R(K) = R(KPIV)
          R(KPIV) = TEMP
180  KPI = K + 1
          DO 190 I=KPI,N
              R(I) = R(I) - D(I,K)*R(K)
190  CONTINUE
200  CONTINUE
C SOLVE U*E=R AND STORE IN R. ALSO LET X=X+E.
      R(N) = R(N)/D(N,N)
      GO TO (210, 220, 210, 220), OPTION
210  X(N) = R(N)
      GO TO 230
220  X(N) = X(N) + R(N)
230  DO 270 II=1,NM1
          I = N - II
          SUM = 0.0D0
          IPI = I + 1
          DO 240 J=IPI,N
              SUM = SUM + D(I,J)*R(J)
240  CONTINUE
          R(I) = (R(I)-SUM)/D(I,I)
          GO TO (250, 260, 250, 260), OPTION
250  X(I) = R(I)
          GO TO 270
260  X(I) = X(I) + R(I)
270  CONTINUE
C SOLUTION OF LINEAR SYSTEM IS COMPLETE. RETURN FOR OPTION=1,3.
      GO TO (280, 290, 280, 290), OPTION
280  RETURN
290  GO TO (130, 300), ISWIT
C CALCULATE ERRORS BASED ON CORRECTION.
300  NORMX = 0.0D0
      NORME = 0.0D0
      DO 310 I=1,N
          NORMX = DMAX1(NORMX,DABS(X(I)))
          NORME = DMAX1(NORME,DABS(R(I)))
310  CONTINUE
      ERROR = NORME/NORMX
      RETURN
      END

```

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      SUBROUTINE IEGAUS(KERNEL, RHFCN, A, B, EP, IFLAG, X, T, NT,
          * NUPPER, MUPPER, W, IER)
C THE INTEGRAL EQUATION BEING SOLVED IS
C      B
C      X(S) - INT KERNEL(S,T)*X(T)*DT = RHFCN(S)
C      A
C THE METHOD BEING USED IS BASED ON THE NYSTROM METHOD WITH
C GAUSSIAN QUADRATURE, WITH AN ITERATIVE TECHNIQUE OF SOLUTION
C FOR THE RESULTING LINEAR SYSTEM.
C KERNEL THESE ARE DOUBLE PRECISION FUNCTIONS OF TWO AND ONE
C RHFCN VARIABLES, RESPECTIVELY. THEY MUST BE DECLARED IN AN
C EXTERNAL STATEMENT IN THE PROGRAM CALLING IEGAUS.
C EP THE DESIRED ERROR. THE VARIABLE EP IS CHANGED ON
C COMPLETION OF THE PROGRAM. SEE THE DISCUSSION OF IER

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C          AND IFLAG FOR MORE INFORMATION. IEG 150
C IFLAG =0 EP IS INTERPRETED AS AN ABSOLUTE ERROR TOLERANCE. IEG 160
C          =1 EP IS INTERPRETED AS A RELATIVE ERROR TOLERANCE. IEG 170
C X        THE COMPUTED APPROXIMATE SOLUTION OF THE INTEGRAL IEG 180
C          EQUATION, EVALUATED AT THE NODE POINTS IN T, IS IEG 190
C          STORED IN X ON COMPLETION OF THE ROUTINE. THIS IS IEG 200
C          TRUE IRREGARDLESS OF WHETHER OR NOT THE DESIRED ERROR IEG 210
C          TOLERANCE WAS ATTAINED. IEG 220
C T        CONTAINS THE NODE POINTS AT WHICH THE SOLUTION OF THE IEG 230
C          INTEGRAL EQUATION IS DESIRED. SEE THE VARIABLE NT FOR IEG 240
C          MORE INFORMATION. IEG 250
C NT       IF NT=0, THEN T AND X WILL BE SET EQUAL TO THE FINAL IEG 260
C          GAUSSIAN NODES AND THE CORRESPONDING SOLUTION VALUES, IEG 270
C          AND NT WILL BE SET TO THE NUMBER OF THE SOLUTION IEG 280
C          VALUES STORED IN X AND T. THE ARRAYS T AND X SHOULD IEG 290
C          HAVE DIMENSION AT LEAST MUPPER, ASSIGNED IN THE IEG 300
C          CALLING PROGRAM. IEG 310
C          IF NT .GT. 0, THEN T CONTAINS NT USER SUPPLIED NODES IEG 320
C          AT WHICH THE SOLUTION X IS DESIRED. IEG 330
C MUPPER   AN UPPER LIMIT ON THE VARIABLE N IN THIS PROGRAM. IEG 340
C          N IS THE ORDER OF A LINEAR SYSTEM WHICH IS BEING IEG 350
C          USED TO ITERATIVELY SOLVE A LARGER LINEAR SYSTEM OF IEG 360
C          ORDER M WHICH APPROXIMATES THE ABOVE INTEGRAL IEG 370
C          EQUATION. IEG 380
C MUPPER   AN UPPER LIMIT ON THE VARIABLE M IN THE PROGRAM. IEG 390
C          N AND M ARE ALWAYS POWERS OF TWO. IEG 400
C W        TEMPORARY WORKING STORAGE FOR THE PROGRAM. IT MUST IEG 410
C          CONTAIN AT LEAST 5*NU*NU+9*(NU+MU) POSITIONS, WITH IEG 420
C          NU=MUPPER, MU=MUPPER. IEG 430
C IER =0   THIS ERROR COMPLETION CODE MEANS THE ROUTINE WAS IEG 440
C          COMPLETED SATISFACTORILY. EP CONTAINS THE PREDICTED IEG 450
C          ERROR. IEG 460
C          =1 THE ERROR TEST WAS NOT SATISFIED. EP CONTAINS THE IEG 470
C          PREDICTED ERROR. IEG 480
C          =2 THE ERROR TEST WAS NOT SATISFIED. EP HAS BEEN SET IEG 490
C          TO ZERO. IEG 500
C          =3 THE ORIGINAL VALUE OF EP WAS TOO SMALL, DUE TO IEG 510
C          POSSIBLE ILL-CONDITIONING PROBLEMS IN THE INTEGRAL IEG 520
C          EQUATION. THE VALUE OF EP WAS RESET TO A MORE IEG 530
C          REALISTIC VALUE, AND THAT TOLERANCE WAS ATTAINED. IEG 540
C          =4 THE ERROR WAS SATISFACTORY AT THE GAUSSIAN NODE IEG 550
C          POINTS (IER=0), BUT THE INTERPOLATION PROCESS(DUE TO IEG 560
C          NT .GT. 0) MAY NOT PRESERVE THIS ACCURACY. CHECK THE IEG 570
C          VALUE OF NORM(K) FOR POSSIBLE INDICATIONS THAT THE IEG 580
C          INTEGRAL EQUATION MAY BE ALMOST FIRST KIND. SUCH IEG 590
C          EQUATIONS ARE QUITE ILL-CONDITIONED. THE ERROR IN EP IEG 600
C          IS THE PREDICTED ERROR FOR THE SOLUTION AT THE IEG 610
C          GAUSSIAN NODE POINTS OF ORDER MFINAL. IEG 620
C          =5 THE ANALOGUE OF IER=4, BUT WITH IER=1 AT THE IEG 630
C          GAUSSIAN NODE POINTS. IEG 640
C          =6 THE ANALOGUE OF IER=4, BUT WITH IER=3 AT THE IEG 650
C          GAUSSIAN NODE POINTS. IEG 660
C IEGAUS IS PRESENTLY LIMITED TO MUPPER .LE. 100. THIS IS ENTIRELY IEG 670
C DUE TO A RESTRICTION IN THE PROGRAM LINSYS. TO REMOVE THIS IEG 680
C RESTRICTION, CHANGE THE DIMENSION STATEMENTS IN THE IEG 690
C COMMON/LINEAR/ STATEMENTS IN LINSYS AND THE SUBPROGRAM IEGB. IEG 700
C          DOUBLE PRECISION KERNEL, RHFCN, A, B, EP, X, T, W, CUTOFF, IEG 710
C          * ROOTRT, UNITRD, R1, R2, FINLEP, NORMK IEG 720
C          DIMENSION X(1), T(1), W(1) IEG 730
C          EXTERNAL KERNEL, RHFCN IEG 740
C ***** IEG 750
C          * IEG 760
C          COMMON /INFO/ R1, R2, FINLEP, NORMK, NFINAL, MFINAL IEG 770
C          * IEG 780
C          THE NUMBERS IN INFO GIVE ADDITIONAL INFORMATION ABOUT THE * IEG 790
C          FUNCTIONING OF IEGAUS. R1 IS THE ITERATIVE RATE OF CONVERGENCE * IEG 800
C          IN THE MOST RECENTLY COMPUTED LINEAR SYSTEM. R2 IS THE RATE OF * IEG 810
C          CONVERGENCE OF THE GAUSSIAN QUADRATURE VARIANT OF THE NYSTROM * IEG 820
C          METHOD. FINLEP IS THE FINAL VALUE OF EP USED AS THE DESIRED * IEG 830
C          ERROR TOLERANCE. USUALLY FINLEP WILL EQUAL THE INPUT VALUE OF * IEG 840
C          EP, UNLESS EP WAS MUCH TOO SMALL. NORMK IS AN APPROXIMATE * IEG 850
C          VALUE FOR THE NORM OF THE INTEGRAL OPERATOR K, AND IT IS * IEG 860
C          CALCULATED ONLY IF NT .GT. 0. * IEG 870
C          NFINAL AND MFINAL ARE THE FINAL VALUES OF N AND M USED IN * IEG 880
C          IEGAUS. IF NFINAL=MFINAL, THEN ITERATION WAS NOT INVOKED * IEG 890

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      DO 140 J=1,N
      SUM = SUM + A(I,J)*X(J)
140  CONTINUE
      R(I) = B(I) - SUM
150  CONTINUE
      NORMB = 0.0D0
      NORMR = 0.0D0
      DO 160 I=1,N
      NORMB = DMAX1(NORMB,DABS(B(I)))
      NORMR = DMAX1(NORMR,DABS(R(I)))
160  CONTINUE
      RELRSD = NORMR/NORMB
      ISWIT = 2
C SOLVE P*L*Z=R AND STORE IN R.
170  NM1 = N - 1
      DO 200 K=1,NM1
      IF (PIVOT(K).EQ.K) GO TO 180
      KPIV = PIVOT(K)
      TEMP = R(K)
      R(K) = R(KPIV)
      R(KPIV) = TEMP
180  KPI = K + 1
      DO 190 I=KPI,N
      R(I) = R(I) - D(I,K)*R(K)
190  CONTINUE
200  CONTINUE
C SOLVE U*E=R AND STORE IN R. ALSO LET X=X+E.
      R(N) = R(N)/D(N,N)
      GO TO (210, 220, 210, 220), OPTION
210  X(N) = R(N)
      GO TO 230
220  X(N) = X(N) + R(N)
230  DO 270 II=1,NM1
      I = N - II
      SUM = 0.0D0
      IP1 = I + 1
      DO 240 J=IP1,N
      SUM = SUM + D(I,J)*R(J)
240  CONTINUE
      R(I) = (R(I)-SUM)/D(I,I)
      GO TO (250, 260, 250, 260), OPTION
250  X(I) = R(I)
      GO TO 270
260  X(I) = X(I) + R(I)
270  CONTINUE
C SOLUTION OF LINEAR SYSTEM IS COMPLETE. RETURN FOR OPTION=1,3.
      GO TO (280, 290, 280, 290), OPTION
280  RETURN
290  GO TO (130, 300), ISWIT
C CALCULATE ERRORS BASED ON CORRECTION.
300  NORMX = 0.0D0
      NORME = 0.0D0
      DO 310 I=1,N
      NORMX = DMAX1(NORMX,DABS(X(I)))
      NORME = DMAX1(NORME,DABS(R(I)))
310  CONTINUE
      ERROR = NORME/NORMX
      RETURN
      END

```

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      SUBROUTINE IEGAUS(KERNEL, RHFCN, A, B, EP, IFLAG, X, T, NT,
      * NUPPER, MUPPER, W, IER)
C THE INTEGRAL EQUATION BEING SOLVED IS
C      B
C      X(S) - INT KERNEL(S,T)*X(T)*DT = RHFCN(S)
C      A
C THE METHOD BEING USED IS BASED ON THE NYSTROM METHOD WITH
C GAUSSIAN QUADRATURE, WITH AN ITERATIVE TECHNIQUE OF SOLUTION
C FOR THE RESULTING LINEAR SYSTEM.
C KERNEL THESE ARE DOUBLE PRECISION FUNCTIONS OF TWO AND ONE
C RHFCN VARIABLES, RESPECTIVELY. THEY MUST BE DECLARED IN AN
C EXTERNAL STATEMENT IN THE PROGRAM CALLING IEGAUS.
C EP THE DESIRED ERROR. THE VARIABLE EP IS CHANGED ON
C COMPLETION OF THE PROGRAM. SEE THE DISCUSSION OF IER

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C          AND IFLAG FOR MORE INFORMATION. IEG 150
C IFLAG =0 EP IS INTERPRETED AS AN ABSOLUTE ERROR TOLERANCE. IEG 160
C          =1 EP IS INTERPRETED AS A RELATIVE ERROR TOLERANCE. IEG 170
C X        THE COMPUTED APPROXIMATE SOLUTION OF THE INTEGRAL IEG 180
C          EQUATION, EVALUATED AT THE NODE POINTS IN T, IS IEG 190
C          STORED IN X ON COMPLETION OF THE ROUTINE. THIS IS IEG 200
C          TRUE IRREGARDLESS OF WHETHER OR NOT THE DESIRED ERROR IEG 210
C          TOLERANCE WAS ATTAINED. IEG 220
C T        CONTAINS THE NODE POINTS AT WHICH THE SOLUTION OF THE IEG 230
C          INTEGRAL EQUATION IS DESIRED. SEE THE VARIABLE NT FOR IEG 240
1C        MORE INFORMATION. IEG 250
C NT       IF NT=0, THEN T AND X WILL BE SET EQUAL TO THE FINAL IEG 260
C          GAUSSIAN NODES AND THE CORRESPONDING SOLUTION VALUES, IEG 270
C          AND NT WILL BE SET TO THE NUMBER OF THE SOLUTION IEG 280
C          VALUES STORED IN X AND T. THE ARRAYS T AND X SHOULD IEG 290
C          HAVE DIMENSION AT LEAST MUPPER, ASSIGNED IN THE IEG 300
C          CALLING PROGRAM. IEG 310
C          IF NT .GT. 0, THEN T CONTAINS NT USER SUPPLIED NODES IEG 320
C          AT WHICH THE SOLUTION X IS DESIRED. IEG 330
C NUPPER   AN UPPER LIMIT ON THE VARIABLE N IN THIS PROGRAM. IEG 340
C          N IS THE ORDER OF A LINEAR SYSTEM WHICH IS BEING IEG 350
C          USED TO ITERATIVELY SOLVE A LARGER LINEAR SYSTEM OF IEG 360
C          ORDER M WHICH APPROXIMATES THE ABOVE INTEGRAL IEG 370
C          EQUATION. IEG 380
C MUPPER   AN UPPER LIMIT ON THE VARIABLE M IN THE PROGRAM. IEG 390
C          N AND M ARE ALWAYS POWERS OF TWO. IEG 400
C W        TEMPORARY WORKING STORAGE FOR THE PROGRAM. IT MUST IEG 410
C          CONTAIN AT LEAST 5*NU*NU+9*(NU+MU) POSITIONS, WITH IEG 420
C          NU=NUPPER, MU=MUPPER. IEG 430
C IER =0   THIS ERROR COMPLETION CODE MEANS THE ROUTINE WAS IEG 440
C          COMPLETED SATISFACTORILY. EP CONTAINS THE PREDICTED IEG 450
C          ERROR. IEG 460
C          =1 THE ERROR TEST WAS NOT SATISFIED. EP CONTAINS THE IEG 470
C          PREDICTED ERROR. IEG 480
C          =2 THE ERROR TEST WAS NOT SATISFIED. EP HAS BEEN SET IEG 490
C          TO ZERO. IEG 500
C          =3 THE ORIGINAL VALUE OF EP WAS TOO SMALL, DUE TO IEG 510
C          POSSIBLE ILL-CONDITIONING PROBLEMS IN THE INTEGRAL IEG 520
C          EQUATION. THE VALUE OF EP WAS RESET TO A MORE IEG 530
C          REALISTIC VALUE, AND THAT TOLERANCE WAS ATTAINED. IEG 540
C          =4 THE ERROR WAS SATISFACTORY AT THE GAUSSIAN NODE IEG 550
C          POINTS (IER=0), BUT THE INTERPOLATION PROCESS(DUE TO IEG 560
C          NT .GT. 0) MAY NOT PRESERVE THIS ACCURACY. CHECK THE IEG 570
C          VALUE OF NORM(K) FOR POSSIBLE INDICATIONS THAT THE IEG 580
C          INTEGRAL EQUATION MAY BE ALMOST FIRST KIND. SUCH IEG 590
C          EQUATIONS ARE QUITE ILL-CONDITIONED. THE ERROR IN EP IEG 600
C          IS THE PREDICTED ERROR FOR THE SOLUTION AT THE IEG 610
C          GAUSSIAN NODE POINTS OF ORDER MFINAL. IEG 620
C          =5 THE ANALOGUE OF IER=4, BUT WITH IER=1 AT THE IEG 630
C          GAUSSIAN NODE POINTS. IEG 640
C          =6 THE ANALOGUE OF IER=4, BUT WITH IER=3 AT THE IEG 650
C          GAUSSIAN NODE POINTS. IEG 660
C IEGAUS IS PRESENTLY LIMITED TO NUPPER .LE. 100. THIS IS ENTIRELY IEG 670
C DUE TO A RESTRICTION IN THE PROGRAM LINSYS. TO REMOVE THIS IEG 680
C RESTRICTION, CHANGE THE DIMENSION STATEMENTS IN THE IEG 690
C COMMON/LINEAR/ STATEMENTS IN LINSYS AND THE SUBPROGRAM IEGS. IEG 700
C          DOUBLE PRECISION KERNEL, RHFCN, A, B, EP, X, T, W, CUTOFF, IEG 710
C          * ROOTRT, UNITRD, R1, R2, FINLEP, NORMK IEG 720
C          DIMENSION X(1), T(1), W(1) IEG 730
C          EXTERNAL KERNEL, RHFCN IEG 740
C ***** IEG 750
C          * IEG 760
C          COMMON /INFO/ R1, R2, FINLEP, NORMK, NFINAL, MFINAL IEG 770
C          * IEG 780
C          THE NUMBERS IN INFO GIVE ADDITIONAL INFORMATION ABOUT THE * IEG 790
C          FUNCTIONING OF IEGAUS. R1 IS THE ITERATIVE RATE OF CONVERGENCE * IEG 800
C          IN THE MOST RECENTLY COMPUTED LINEAR SYSTEM. R2 IS THE RATE OF * IEG 810
C          CONVERGENCE OF THE GAUSSIAN QUADRATURE VARIANT OF THE NYSTROM * IEG 820
C          METHOD. FINLEP IS THE FINAL VALUE OF EP USED AS THE DESIRED * IEG 830
C          ERROR TOLERANCE. USUALLY FINLEP WILL EQUAL THE INPUT VALUE OF * IEG 840
C          EP, UNLESS EP WAS MUCH TOO SMALL. NORMK IS AN APPROXIMATE * IEG 850
C          VALUE FOR THE NORM OF THE INTEGRAL OPERATOR K, AND IT IS * IEG 860
C          CALCULATED ONLY IF NT .GT. 0. * IEG 870
C          NFINAL AND MFINAL ARE THE FINAL VALUES OF N AND M USED IN * IEG 880
C          IEGAUS. IF NFINAL=MFINAL, THEN ITERATION WAS NOT INVOKED * IEG 890

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C      SUCCESSFULLY. * IEG 900
C * IEG 910
C ***** IEG 920
C * IEG 930
C      DATA UNITRD /2.22D-16/ IEG 940
C * IEG 950
C      UNITRD IS THE SMALLEST NUMBER U FOR WHICH * IEG 960
C      1+U .GT. 1 . * IEG 970
C      UNITRD VARIES WITH THE COMPUTER AND WITH THE ARITHMETIC BEING * IEG 980
C      USED. TO CHANGE TO ANOTHER ARITHMETIC, UNITRD MUST BE CHANGED. * IEG 990
C      BUT UNITRD ALSO REFLECTS THE ACCURACY OF THE CONSTANTS USED * IEG 1000
C      IN SUBROUTINE WANDT. * IEG 1010
C * IEG 1020
C ***** IEG 1030
C      DATA CUTOFF /0.5D0/, ROOTRT /0.1D0/ IEG 1040
C      SET UP THE RELATIVE BASE ADDRESSES FOR THE VARIOUS ARRAYS INTO IEG 1050
C      WHICH W IS TO BE SPLIT. IEG 1060
C      N = NUPPER IEG 1070
C      M = MUPPER IEG 1080
C      NSQ = N*N IEG 1090
C      I1 = 1 IEG 1100
C      I2 = I1 + NSQ IEG 1110
C      I3 = I2 + NSQ IEG 1120
C      I4 = I3 + NSQ/2 IEG 1130
C      I5 = I4 + NSQ/2 IEG 1140
C      I6 = I5 + M IEG 1150
C      I7 = I6 + M IEG 1160
C      I8 = I7 + N IEG 1170
C      I9 = I8 + N IEG 1180
C      I10 = I9 + M IEG 1190
C      I11 = I10 + N IEG 1200
C      I12 = I11 + M IEG 1210
C      I13 = I12 + M IEG 1220
C      I14 = I13 + M IEG 1230
C      I15 = I14 + N IEG 1240
C      I16 = I15 + M IEG 1250
C      I17 = I16 + M IEG 1260
C      I18 = I17 + N IEG 1270
C      I19 = I18 + M IEG 1280
C      I20 = I19 + 4*N IEG 1290
C      I21 = I20 + NSQ IEG 1300
C      NHALF = N/2 IEG 1310
C      CALL IEKS(KERNEL, RHFCN, A, B, EP, IFLAG, X, T, NT, NUPPER, IEG 1320
* MUPPER, IER, CUTOFF, ROOTRT, UNITRD, NHALF, W(I1), W(I2), IEG 1330
* W(I3), W(I4), W(I5), W(I6), W(I7), W(I8), W(I9), W(I10), IEG 1340
* W(I11), W(I12), W(I13), W(I14), W(I15), W(I16), W(I17), IEG 1350
* W(I18), W(I19), W(I20), W(I21)) IEG 1360
      RETURN IEG 1370
      END IEG 1380

      SUBROUTINE IEKS(KERNEL, RHFCN, A, B, EP, IFLAG, X, T, NT, IEG 10
* NUP, MUP, IER, CUTOFF, ROOTRT, UNITRD, NHALF, LUFAC, KMM, IEG 20
* KMN, KNM, RHS, R, RH, DELN, TM, TN, XM, XMZ, WM, WN, OLDX, IEG 30
* SAVE, XN, SAVE2, ASIDE, ASIDE3, IMKNN) IEG 40
C THIS ROUTINE CONTROLS THE SOLUTION OF THE INTEGRAL EQUATION. IEG 50
      DOUBLE PRECISION KERNEL, RHFCN, A, B, EP, X, T, CUTOFF, IEG 60
* ROOTRT, UNITRD, LUFAC, KMM, KMN, KNM, RHS, R, RH, DELN, TM, IEG 70
* TN, XM, XMZ, WM, WN, OLDX, SAVE, XN, SAVE2, ASIDE, ASIDE2, IEG 80
* ASIDE3, IMKNN, RESID, SCALE, ELINSY, RELRSD, XNORM, PASTC, IEG 90
* PASTRE, R1, R2, FINLEP, NORMK, DFLOAT, NORM, CONEW, RMIN, IEG 100
* RIRAT, COND, AVERR, EPS, DET, RELMIN, DMIN1, DMAX1, DSQRT, IEG 110
* ERROR, NUMR2, DENR2, TEMP2, RELERR, NUMR1, DENR1, RATE, IEG 120
* TEMP, RT, ESTERR, TEST IEG 130
      INTEGER FLAG, OLDM IEG 140
      DIMENSION X(1), T(1), LUFAC(NUP,NUP), KMM(NUP,NUP), IEG 150
* RHS(MUP), KNM(NHALF,NUP), KMN(NUP,NHALF), R(MUP), RH(NUP), IEG 160
* DELN(NUP), TM(MUP), TN(NUP), XM(MUP), XMZ(MUP), WM(MUP), IEG 170
* WN(NUP), OLDX(MUP), SAVE(MUP), XN(NUP), SAVE2(MUP), IEG 180
* ASIDE(NUP,4), ASIDE2(5), ASIDE3(NUP,NUP), IMKNN(NUP,NUP) IEG 190
      COMMON /LINEAR/ RESID(100), SCALE(100), ELINSY, RELRSD, IEG 200
* IPIVOT(100) IEG 210
      COMMON /INFO/ R1, R2, FINLEP, NORMK, NFINAL, MFINAL IEG 220
      EXTERNAL KERNEL, RHFCN IEG 230
C INITIALIZATION IEG 240

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      LOOP = 1
      N = 2
      R2 = 0.5D0
      M = 2*N
      RIRAT = ROOTRT
      COND = 1.0D0
      PASTC = 1.0D0
      PASTRE = 0.0D0
      EPS = EP
C STAGE A. DIRECT SOLUTION OF LINEAR SYSTEM (I-KN)*XN=RHS, WHILE
C TRYING TO FIND A GOOD APPROXIMATE INVERSE TO IMPLEMENT
C ITERATIVE METHOD OF SOLUTION.
C CREATE THE NODES AND WEIGHTS TN(I) AND WN(I), I=1,...,N
      CALL WANDT(WN, TN, N, A, B)
C SET UP MATRIX FOR (I-KN)*XN=RHFCN
      DO 20 J=1,N
        DO 10 I=1,N
          IMKNN(I,J) = -WN(J)*KERNEL(TN(I),TN(J))
10      CONTINUE
          XMZ(J) = RHFCN(TN(J))
          IMKNN(J,J) = IMKNN(J,J) + 1.0D0
20      CONTINUE
      GO TO 60
C THIS IS ENTRANCE FOR AN INCREASED VALUE OF N, USING PREVIOUSLY
C STORED VALUES IN KMM TO DEFINE MATRIX FOR (I-KN)*XN=RHFCN WITH
C NEW VALUE OF N.
      30 DO 50 J=1,N
        DO 40 I=1,N
          IMKNN(I,J) = -KMM(I,J)
40      CONTINUE
          WN(J) = WM(J)
          TN(J) = TM(J)
          XMZ(J) = RHS(J)
          IMKNN(J,J) = IMKNN(J,J) + 1.0D0
50      CONTINUE
C THIS IS THE ENTRANCE WHEN ITERATION IN STAGE B FAILS AND WE NEED
C TO INCREASE N TO OBTAIN A BETTER ITERATIVE RATE.
      60 CONTINUE
C SOLVE (I-KN)*XN=RHFCN AT ALL TN(I).ALSO OBTAIN THE LU
C DECOMPOSITION FOR LATER USE IN THE STAGE B ITERATIVE METHOD.
C *****
C CALL LINSYS(IMKNN, LUFAC, N, XMZ, XN, 2, DET, NUP)
C
C LINSYS IS A GENERAL LINEAR EQUATION SOLVER. IT HAS SPECIAL
C OPTIONS WHICH ARE USED IN THE FOLLOWING PROGRAM, AND THUS
C LINSYS SHOULD NOT BE REPLACED BY ANOTHER LINEAR EQUATION
C PROGRAM. LINSYS IS ALSO USED IN THE SUBROUTINE ITERT.
C *****
C COND = CONEW(COND,ELINSY,RELRSD,AVERR,PASTC,PASTRE)
      RELMIN = RMIN(N,N,COND,UNITRD,AVERR)
      IF (LOOP.EQ.1) GO TO 100
      IF (LOOP.EQ.2) GO TO 80
C SET UP APPROXIMATE RATE OF CONVERGENCE OF SOLUTIONS XN TO TRUE
C SOLUTION X. ALSO SET UP DESIRED RATIO FOR ITERATIVE METHOD.
      NUMR2 = NORM(XN,OLDX,N,1)
      R2 = DMIN1(0.5D0,DMAX1(NUMR2/DENR2,1.0D-4))
      RIRAT = DMIN1(ROOTRT,DSQRT(R2))
C CHECK FOR ERROR IN XN USING TEST INVOLVING R2 AND OLDX,ACCORDING
C TO THEORY FOR ASYMPTOTIC ERROR BOUNDS. MODIFY ERROR IF IT IS
C OUTSIDE PRECISION RANGE OF COMPUTER, POSSIBLY DUE TO
C ILL-CONDITIONING.
      70 ERROR = (R2/(1.0D0-R2))*NUMR2
      XNORM = NORM(XN,XN,N,0)
      RELERR = ERROR/XNORM
      IF (IFLAG.EQ.0) EPS = DMAX1(EP,XNORM*RELMIN)
      IF (IFLAG.EQ.1) EPS = DMAX1(EP,RELMIN)
      IF (IFLAG.EQ.1) ERROR = DMAX1(RELERR,RELMIN)
      IF ((IFLAG.EQ.0) .AND. (RELERR.LT.RELMIN)) ERROR =
      * RELMIN*XNORM
      IF (ERROR.LE.EPS) GO TO 90
      DENR2 = NUMR2
      GO TO 100
C ENTRANCE FOR LOOP=2.
      80 NUMR2 = NORM(XN,OLDX,N,1)

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IEG 250
IEG 260
IEG 270
IEG 280
IEG 290
IEG 300
IEG 310
IEG 320
IEG 330
IEG 340
IEG 350
IEG 360
IEG 370
IEG 380
IEG 390
IEG 400
IEG 410
IEG 420
IEG 430
IEG 440
IEG 450
IEG 460
IEG 470
IEG 480
IEG 490
IEG 500
IEG 510
IEG 520
IEG 530
IEG 540
IEG 550
IEG 560
IEG 570
IEG 580
IEG 590
IEG 600
IEG 610
IEG 620
IEG 630
IEG 640
IEG 650
IEG 660
IEG 670
IEG 680
IEG 690
IEG 700
IEG 710
IEG 720
IEG 730
IEG 740
IEG 750
IEG 760
IEG 770
IEG 780
IEG 790
IEG 800
IEG 810
IEG 820
IEG 830
IEG 840
IEG 850
IEG 860
IEG 870
IEG 880
IEG 890
IEG 900
IEG 910
IEG 920
IEG 930
IEG 940
IEG 950
IEG 960
IEG 970
IEG 980
IEG 990
IEG 1000

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DENR2 = 0.0D0
GO TO 70
C EXIT FOR SUCCESSFUL RETURN. ITERATION WAS NOT NECESSARY.
90 CALL LEAVE(0, N, N, XN, TN, WN, ERROR, KERNEL, RHFCN, EP,
* IFLAG, X, T, NT, IER, EPS, ELINSY, TN, WN, TM, WM, XM, XMZ,
* KMM, KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, NUP, NHALF,
* XNORM)
RETURN
C ATTEMPT TO SOLVE (I-KM)*XM=RHFCN ITERATIVELY, CHECKING TO SEE IF
C THE RATE OF CONVERGENCE IS SUFFICIENTLY FAST SO AS TO ENTER
C STAGE B.
C CALCULATE TM(I) AND WM(I), I=1,...,M.
100 CALL WANDT(WM, TM, M, A, B)
FLAG = 0
C CALCULATE INITIAL GUESS XMZ FOR ITERATION METHOD.
CALL INTERP(TM, WM, XMZ, M, TN, WN, XN, N, KERNEL, RHFCN,
* RHS, KMN, NHALF, NUP)
DO 110 I=1,M
OLDX(I) = XMZ(I)
110 CONTINUE
C CALCULATE FIRST ITERATE.
CALL ITERT(KERNEL, RHFCN, N, TN, WN, M, TM, WM, XM, XMZ, KMM,
* KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, NUP, NHALF, FLAG)
COND = CONEW(COND,ELINSY,RELRSD,AVERR,PASTC,PASTRE)
DENR1 = NORM(XM, XMZ, M, 1)
FLAG = 1
DO 120 I=1,M
XMZ(I) = XM(I)
120 CONTINUE
C CALCULATE SECOND ITERATE.
CALL ITERT(KERNEL, RHFCN, N, TN, WN, M, TM, WM, XM, XMZ, KMM,
* KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, NUP, NHALF, FLAG)
COND = CONEW(COND,ELINSY,RELRSD,AVERR,PASTC,PASTRE)
NUMR1 = NORM(XM, XMZ, M, 1)
C CHECK ON THE SPEED OF CONVERGENCE OF ITERATIVE METHOD. IF IT IS
C SUFFICIENTLY RAPID, THEN FIX N AND GO TO STAGE B.
R1 = NUMR1/DENR1
RATE = R1
IF (M.GT.NUP) GO TO 170
IF (R1.LE.R1RAT) GO TO 130
C THE ITERATION DID NOT WORK WELL ENOUGH, AND STAGE A IS TO BE
C REPEATED. RE-INITIALIZE FOR SOLVING (I-KN)*XN=RHFCN AGAIN
C WITH A LARGER N.
N = M
LOOP = LOOP + 1
M = 2*N
GO TO 30
C THE ITERATIVE RATE IS SUFFICIENTLY RAPID, AND CONTROL WILL GO TO
C STAGE B. SAVE INFORMATION IN CASE STAGE B ABORTS AT A LARGER
C VALUE OF M AND STAGE A HAS TO BE RETURNED TO.
130 DO 140 I=1,M
ASIDE(I,1) = OLDX(I)
ASIDE(I,2) = WM(I)
ASIDE(I,3) = TM(I)
ASIDE(I,4) = RHS(I)
140 CONTINUE
ASIDE2(1) = LOOP
ASIDE2(3) = R2
ASIDE2(4) = DENR2
ASIDE2(5) = R1RAT
DO 160 J=1,M
DO 150 I=1,M
ASIDE3(I,J) = KMM(I,J)
150 CONTINUE
160 CONTINUE
C STAGE B. ITERATIVE METHOD OF SOLUTION OF (I-KM)*XM=RHS.
170 OLDM = N
ASIDE2(2) = M
IF (R1.LE.CUTOFF) GO TO 200
C THE ITERATES ARE CONVERGING VERY SLOWLY OR NOT AT ALL. THUS
C RETURN WITHOUT FURTHER ATTEMPTS TO LESSEN THE ERROR.
IF (LOOP.NE.1) GO TO 190
180 CALL LEAVE(2, N, N, XN, TN, WN, 0.0D0, KERNEL, RHFCN, EP,
* IFLAG, X, T, NT, IER, EPS, ELINSY, TN, WN, TM, WM, XM, XMZ,
* KMM, KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, NUP, NHALF,

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IEG 1010
IEG 1020
IEG 1030
IEG 1040
IEG 1050
IEG 1060
IEG 1070
IEG 1080
IEG 1090
IEG 1100
IEG 1110
IEG 1120
IEG 1130
IEG 1140
IEG 1150
IEG 1160
IEG 1170
IEG 1180
IEG 1190
IEG 1200
IEG 1210
IEG 1220
IEG 1230
IEG 1240
IEG 1250
IEG 1260
IEG 1270
IEG 1280
IEG 1290
IEG 1300
IEG 1310
IEG 1320
IEG 1330
IEG 1340
IEG 1350
IEG 1360
IEG 1370
IEG 1380
IEG 1390
IEG 1400
IEG 1410
IEG 1420
IEG 1430
IEG 1440
IEG 1450
IEG 1460
IEG 1470
IEG 1480
IEG 1490
IEG 1500
IEG 1510
IEG 1520
IEG 1530
IEG 1540
IEG 1550
IEG 1560
IEG 1570
IEG 1580
IEG 1590
IEG 1600
IEG 1610
IEG 1620
IEG 1630
IEG 1640
IEG 1650
IEG 1660
IEG 1670
IEG 1680
IEG 1690
IEG 1700
IEG 1710
IEG 1720
IEG 1730
IEG 1740
IEG 1750

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* XNORM) IEG 1760
RETURN IEG 1770
190 CALL LEAVE(1, N, N, XN, TN, WN, ERROR, KERNEL, RHFCN, EP, IEG 1780
* IFLAG, X, T, NT, IER, EPS, ELINSY, TN, WN, TM, WM, XM, XMZ, IEG 1790
* KMM, KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, NUP, NHALF, IEG 1800
* XNORM) IEG 1810
RETURN IEG 1820
C TEST TO SEE IF THE CURRENT ITERATE XM IS SUFFICIENTLY ACCURATE IEG 1830
C COMPARED TO THE TRUE XM. IEG 1840
200 RATE = R1*RATE IEG 1850
TEMP = NORM(XM,OLDX,M,1) IEG 1860
IF (LOOP.EQ.1) TEMP2 = 0.5D0 IEG 1870
IF (LOOP.GT.1) TEMP2 = TEMP/DENR2 IEG 1880
RT = DMIN1(0.01D0,DMAX1(TEMP2,0.0001D0))/2.0D0 IEG 1890
XNORM = NORM(XM,XM,M,0) IEG 1900
ESTERR = (RT/(1.0D0-RT))*TEMP/XNORM IEG 1910
IF (ESTERR.LT.RELMIN) ESTERR = RELMIN IEG 1920
ESTERR = ESTERR*XNORM IEG 1930
TEST = ((1.0D0-R1)/R1)*ESTERR IEG 1940
IF (NUMR1.LE.TEST) GO TO 260 IEG 1950
C ITERATE NOT SUFFICIENTLY ACCURATE. INITIALIZE FOR COMPUTATION IEG 1960
C OF ANOTHER ITERATE. IEG 1970
DENR1 = NUMR1 IEG 1980
DO 210 I=1,M IEG 1990
XMZ(I) = XM(I) IEG 2000
210 CONTINUE IEG 2010
CALL ITERT(KERNEL, RHFCN, N, TN, WN, M, TM, WM, XM, XMZ, KMM, IEG 2020
* KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, NUP, NHALF, FLAG) IEG 2030
COND = CONEW(COND,ELINSY,RELRS,AVERR,PASTC,PASTRE) IEG 2040
NUMR1 = NORM(XM,XMZ,M,1) IEG 2050
R1 = NUMR1/DENR1 IEG 2060
IF (R1.LE.CUTOFF) GO TO 200 IEG 2070
C THIS IS ENTRANCE FOR CASE WHERE ITERATION FAILS IN STAGE B. IEG 2080
C PARAMETERS MUST BE RESET FOR A RETURN TO STAGE A OR FOR AN IEG 2090
C ABORTIVE EXIT IF N CANNOT BE INCREASED ANY FURTHER. IEG 2100
220 MNEW = ASIDE2(2) IEG 2110
IF (MNEW.GT.NUP) GO TO 250 IEG 2120
N = MNEW IEG 2130
DO 240 J=1,N IEG 2140
DO 230 I=1,N IEG 2150
IMKNN(I,J) = -ASIDE3(I,J) IEG 2160
230 CONTINUE IEG 2170
OLDX(J) = ASIDE(J,1) IEG 2180
WN(J) = ASIDE(J,2) IEG 2190
TN(J) = ASIDE(J,3) IEG 2200
XMZ(J) = ASIDE(J,4) IEG 2210
IMKNN(J,J) = IMKNN(J,J) + 1.0D0 IEG 2220
240 CONTINUE IEG 2230
M = 2*N IEG 2240
LOOP = ASIDE2(1) + 1.0D0 IEG 2250
R2 = ASIDE2(3) IEG 2260
DENR2 = ASIDE2(4) IEG 2270
R1RAT = ASIDE2(5) IEG 2280
GO TO 60 IEG 2290
C ABORTIVE EXIT FROM STAGE B. N CANNOT BE INCREASED FURTHER, AND IEG 2300
C R1 IS NOT SUFFICIENTLY SMALL. IEG 2310
250 IF (LOOP.EQ.1) GO TO 180 IEG 2320
CALL WANDT(WM, TM, OLDM, A, B) IEG 2330
CALL LEAVE(1, N, OLDM, SAVE, TM, WM, ERROR, KERNEL, RHFCN, IEG 2340
* EP, IFLAG, X, T, NT, IER, EPS, ELINSY, TN, WN, TM, WM, XM, IEG 2350
* XMZ, KMM, KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, NUP, IEG 2360
* NHALF, XNORM) IEG 2370
RETURN IEG 2380
C AN ACCURATE VALUE OF XM HAS BEEN OBTAINED. R2 IS TO BE TESTED AS IEG 2390
C TO WHETHER IT SHOULD BE RESET. THEN CHECK ERROR IN XM COMPARED IEG 2400
C WITH THE TRUE SOLUTION X. IEG 2410
260 IF (LOOP.EQ.1) GO TO 290 IEG 2420
NUMR2 = TEMP IEG 2430
R2 = DMAX1(1.0D-4,RATE,DMIN1(NUMR2/DENR2,0.5D0)) IEG 2440
DENR2 = NUMR2 IEG 2450
270 ERROR = (R2/(1.0D0-R2))*TEMP IEG 2460
XNORM = NORM(XM,XM,M,0) IEG 2470
RELERR = ERROR/XNORM IEG 2480
RELMIN = RMIN(N,M,COND,UNITRD,AVERR) IEG 2490
IF (IFLAG.EQ.0) EPS = DMAX1(EP,XNORM*RELMIN) IEG 2500
IF (IFLAG.EQ.1) EPS = DMAX1(EP,RELMIN) IEG 2510

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      IF (IFLAG.EQ.1) ERROR = DMAX1(RELERR,RELMIN) IEG 2520
      IF ((IFLAG.EQ.0) .AND. (RELERR.LT.RELMIN)) ERROR = IEG 2530
      * RELMIN*XNORM IEG 2540
      IF (ERROR.GT.EPS) GO TO 280 IEG 2550
      CALL LEAVE(0, N, M, XM, TM, WM, ERROR, KERNEL, RHFCN, EP, IEG 2560
      * IFLAG, X, T, NT, IER, EPS, ELINSY, TN, WN, TM, WM, XM, XMZ, IEG 2570
      * KMM, KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, NUP, NHALF, IEG 2580
      * XNORM) IEG 2590
      RETURN IEG 2600
280 MNEW = 2*M IEG 2610
      IF (MNEW.LE.MUP) GO TO 300 IEG 2620
C M CANNOT BE INCREASED ANY FURTHER. IEG 2630
      CALL LEAVE(1, N, M, XM, TM, WM, ERROR, KERNEL, RHFCN, EP, IEG 2640
      * IFLAG, X, T, NT, IER, EPS, ELINSY, TN, WN, TM, WM, XM, XMZ, IEG 2650
      * KMM, KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, NUP, NHALF, IEG 2660
      * XNORM) IEG 2670
      RETURN IEG 2680
290 DENR2 = TEMP IEG 2690
      LOOP = 2 IEG 2700
      GO TO 270 IEG 2710
C ERROR NOT SUFFICIENTLY SMALL. M IS INCREASED AND TWO MORE IEG 2720
C MORE ITERATES ARE COMPUTED WITH THE NEW M. IEG 2730
300 OLD M = M IEG 2740
      M = MNEW IEG 2750
      DO 310 I=1,OLD M IEG 2760
          SAVE2(I) = WM(I) IEG 2770
          SAVE(I) = TM(I) IEG 2780
310 CONTINUE IEG 2790
      CALL WANDT(WM, TM, M, A, B) IEG 2800
      FLAG = 0 IEG 2810
      CALL INTERP(TM, WM, XMZ, M, SAVE, SAVE2, XM, OLD M, KERNEL, IEG 2820
      * RHFCN, RHS, KMN, NHALF, NUP) IEG 2830
      DO 320 I=1,OLD M IEG 2840
          SAVE(I) = XM(I) IEG 2850
320 CONTINUE IEG 2860
      DO 330 I=1,M IEG 2870
          OLDX(I) = XMZ(I) IEG 2880
330 CONTINUE IEG 2890
      CALL ITERT(KERNEL, RHFCN, N, TN, WN, M, TM, WM, XM, XMZ, KMM, IEG 2900
      * KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, NUP, NHALF, FLAG) IEG 2910
      COND = CONEW(COND,ELINSY,RELRS,AVERR,PASTC,PASTRE) IEG 2920
      DENR1 = NORM(XM,XMZ,M,1) IEG 2930
      FLAG = 1 IEG 2940
      DO 340 I=1,M IEG 2950
          XMZ(I) = XM(I) IEG 2960
340 CONTINUE IEG 2970
      CALL ITERT(KERNEL, RHFCN, N, TN, WN, M, TM, WM, XM, XMZ, KMM, IEG 2980
      * KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, NUP, NHALF, FLAG) IEG 2990
      COND = CONEW(COND,ELINSY,RELRS,AVERR,PASTC,PASTRE) IEG 3000
      NUMR1 = NORM(XM,XMZ,M,1) IEG 3010
      R1 = NUMR1/DENR1 IEG 3020
      RATE = R1 IEG 3030
      IF (R1.LE.CUTOFF) GO TO 200 IEG 3040
      GO TO 220 IEG 3050
      END IEG 3060

      SUBROUTINE WANDT(WV, TV, N, A, B) WAN 10
C INTEGRATION WEIGHTS AND NODES ARE TO BE CALCULATED AND STORED IN WAN 20
C WV AND TV, RESPECTIVELY. N IS ASSUMED TO BE A POWER OF TWO. IF WAN 30
C 2 .LE. N .LE. 256, THEN GAUSSIAN QUADRATURE IS USED. IF N .GT. WAN 40
C 256, THEN THE INTERVAL (A,B) IS DIVIDED N/256 TIMES AND THE 256 WAN 50
C POINT FORMULA IS APPLIED TO EACH SUBINTERVAL. WAN 60
      DOUBLE PRECISION WV, TV, A, B, W, T, FLOOP, H, SCALE, AL, BL, WAN 70
      * S, R, FL WAN 80
      DIMENSION WV(N), TV(N) WAN 90
      DIMENSION W(255), T(255) WAN 100
      DATA T(1), T(2), T(3), T(4), T(5), T(6), T(7), T(8), T(9), WAN 110
      * T(10), T(11), T(12), T(13), T(14), T(15) WAN 120
      * /.577350269189626D0,.861136311594053D0,.339981043584856D0, WAN 130
      * .960289856497536D0,.796666477413627D0,.525532409916329D0, WAN 140
      * .183434642495650D0,.989400934991650D0,.944575023073233D0, WAN 150
      * .865631202387832D0,.755404408355003D0,.617876244402644D0, WAN 160
      * .458016777657227D0,.281603550779259D0,.950125098376374D-1/ WAN 170
      DATA T(16), T(17), T(18), T(19), T(20), T(21), T(22), T(23), WAN 180

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* T(24), T(25), T(26), T(27), T(28), T(29), T(30), T(31)	WAN	190
* /.997263861849482D0, .985611511545268D0, .96476225587506D0,	WAN	200
* .934906075937740D0, .896321155766052D0, .849367613732570D0,	WAN	210
* .794483795967942D0, .732182118740290D0, .663044266930215D0,	WAN	220
* .587715757240762D0, .506899908932229D0, .421351276130635D0,	WAN	230
* .331868602282128D0, .239287362252137D0, .144471961582796D0,	WAN	240
* .483076656877383D-1/ DATA T(32), T(33), T(34), T(35), T(36), T(37), T(38), T(39),	WAN	250
* T(40), T(41), T(42), T(43), T(44), T(45), T(46), T(47)	WAN	260
* /.999305041735772D0, .996340116771955D0, .991013371476744D0,	WAN	270
* .983336253884626D0, .973326827789911D0, .961008799652054D0,	WAN	280
* .946411374858403D0, .929569172131940D0, .910522137078503D0,	WAN	290
* .889315445995114D0, .865999398154093D0, .840629296252580D0,	WAN	300
* .813265315122798D0, .783972358943341D0, .752819907260532D0,	WAN	310
* .719881850171611D0/ DATA T(48), T(49), T(50), T(51), T(52), T(53), T(54), T(55),	WAN	320
* T(56), T(57), T(58), T(59), T(60), T(61), T(62), T(63)	WAN	330
* /.685236313054233D0, .648965471254657D0, .611155355172393D0,	WAN	340
* .571895646202634D0, .531279464019895D0, .489403145707053D0,	WAN	350
* .446366017253464D0, .402270157963992D0, .357220158337668D0,	WAN	360
* .311322871990211D0, .264687162208767D0, .217423643740007D0,	WAN	370
* .169644420423993D0, .121462819296121D0, .729931217877990D-1,	WAN	380
* .243502926634244D-1/ DATA T(64), T(65), T(66), T(67), T(68), T(69), T(70), T(71),	WAN	390
* T(72), T(73), T(74), T(75), T(76), T(77), T(78), T(79)	WAN	400
* /.999824887947132D0, .999077459977376D0, .997733248625514D0,	WAN	410
* .995792758534981D0, .993257112900213D0, .990127818491734D0,	WAN	420
* .986406742724586D0, .982096108435719D0, .977198491463907D0,	WAN	430
* .971716818747137D0, .965654366431965D0, .959014757853700D0,	WAN	440
* .951801961341264D0, .944020287830220D0, .935674388277916D0,	WAN	450
* .926769250878948D0/ DATA T(80), T(81), T(82), T(83), T(84), T(85), T(86), T(87),	WAN	460
* T(88), T(89), T(90), T(91), T(92), T(93), T(94), T(95),	WAN	470
* T(96), T(97) /.917310198080961D0, .907302883401757D0,	WAN	480
* .896753288049158D0, .885667717345397D0, .874052796958032D0,	WAN	490
* .861915468939548D0, .849262987577969D0, .836102915060907D0,	WAN	500
* .822443116955644D0, .808291757507914D0, .793657294762193D0,	WAN	510
* .778548475506412D0, .762974330044095D0, .746944166797062D0,	WAN	520
* .730467566741909D0, .713554377683587D0, .696214708369514D0,	WAN	530
* .678458922447719D0/ DATA T(98), T(99), T(100), T(101), T(102), T(103), T(104),	WAN	540
* T(105), T(106), T(107), T(108), T(109), T(110), T(111),	WAN	550
* T(112), T(113), T(114) /.660297632272646D0,	WAN	560
* .641741692562308D0, .622802193910585D0, .603490456158549D0,	WAN	570
* .583818021628763D0, .563796648226618D0, .543438302412810D0,	WAN	580
* .522755152051175D0, .501759559136144D0, .480464072404172D0,	WAN	590
* .458881419833552D0, .437024501037104D0, .414906379552275D0,	WAN	600
* .392540275033267D0, .369939555349859D0, .347117728597636D0,	WAN	610
* .324088435024413D0/ DATA T(115), T(116), T(117), T(118), T(119), T(120), T(121),	WAN	620
* T(122), T(123), T(124), T(125), T(126), T(127)	WAN	630
* /.300865438877677D0, .277462620177904D0, .253893966422694D0,	WAN	640
* .23017356422660D0, .206315590902079D0, .182334305985337D0,	WAN	650
* .158244042714225D0, .134059199461188D0, .109794231127644D0,	WAN	660
* .854636405045155D-1, .610819696041396D-1, .366637909687335D-1,	WAN	670
* .122236989606158D-1/ DATA T(128), T(129), T(130), T(131), T(132), T(133), T(134),	WAN	680
* T(135), T(136), T(137), T(138), T(139), T(140), T(141),	WAN	690
* T(142) /.999956050018992D0, .999768437409263D0,	WAN	700
* .999430937466261D0, .998943525843409D0, .998306266473006D0,	WAN	710
* .997519252756721D0, .996582602023382D0, .995496454481096D0,	WAN	720
* .994260972922410D0, .992876342608822D0, .991342771207583D0,	WAN	730
* .989660488745065D0, .987829747564861D0, .985850822286126D0,	WAN	740
* .983724009760315D0/ DATA T(143), T(144), T(145), T(146), T(147), T(148), T(149),	WAN	750
* T(150), T(151), T(152), T(153), T(154), T(155), T(156),	WAN	760
* T(157) /.981449629025464D0, .979028021257622D0,	WAN	770
* .976459549719234D0, .973744599704370D0, .970883578480743D0,	WAN	780
* .967876915228489D0, .964725060975706D0, .961428488530732D0,	WAN	790
* .957987692411178D0, .954403188769716D0, .950675515316628D0,	WAN	800
* .946805231239127D0, .942792917117462D0, .938639174837814D0,	WAN	810
* .934344627502003D0/ DATA T(158), T(159), T(160), T(161), T(162), T(163), T(164),	WAN	820
* T(165), T(166), T(167), T(168), T(169), T(170), T(171),	WAN	830
* T(172) /.929909919334006D0, .925335715583316D0,	WAN	840
* .920622702425146D0, .915771586857490D0, .910783096595065D0,	WAN	850
	WAN	860
	WAN	870
	WAN	880
	WAN	890
	WAN	900
	WAN	910
	WAN	920
	WAN	930
	WAN	940

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* .905657979960145D0,.900397005770304D0,.895000963223085D0,      WAN 950
* .889470661777611D0,.883806931033158D0,.878010620604707D0,      WAN 960
* .872082599995488D0,.866023758466555D0,.859835004903376D0,      WAN 970
* .853517267679503D0/                                             WAN 980
DATA T(173), T(174), T(175), T(176), T(177), T(178), T(179),      WAN 990
* T(180), T(181), T(182), T(183), T(184), T(185), T(186),      WAN 1000
* T(187) /.847071494517296D0,.840498652345763D0,                 WAN 1010
* .833799727155505D0,.826975723850813D0,.820027666098917D0,      WAN 1020
* .812956596176432D0,.805763574812999D0,.798449681032171D0,      WAN 1030
* .791016011989546D0,.783463682808184D0,.775793826411326D0,      WAN 1040
* .768007593352446D0,.760106151642655D0,.7520900686575492D0,      WAN 1050
* .743962400549112D0/                                             WAN 1060
DATA T(188), T(189), T(190), T(191), T(192), T(193), T(194),      WAN 1070
* T(195), T(196), T(197), T(198), T(199), T(200), T(201),      WAN 1080
* T(202) /.735722512885918D0,.727372259649652D0,                 WAN 1090
* .718912893459971D0,.710345683304543D0,.701671914348685D0,      WAN 1100
* .692892887742577D0,.684009920426076D0,.675024344931163D0,      WAN 1110
* .665937509182049D0,.656750776292973D0,.647465524363725D0,      WAN 1120
* .638083146272911D0,.628605049469015D0,.619032655759261D0,      WAN 1130
* .609367401096334D0/                                             WAN 1140
DATA T(203), T(204), T(205), T(206), T(207), T(208), T(209),      WAN 1150
* T(210), T(211), T(212), T(213), T(214), T(215), T(216),      WAN 1160
* T(217) /.599610735362968D0,.589764122154454D0,                 WAN 1170
* .579829038559083D0,.569806974936569D0,.559699434694481D0,      WAN 1180
* .549507934062719D0,.539234001866059D0,.528879179294822D0,      WAN 1190
* .518445019673674D0,.507933088228616D0,.497344961852181D0,      WAN 1200
* .486682228866890D0,.475946488786983D0,.465139352078479D0,      WAN 1210
* .454262439917590D0/                                             WAN 1220
DATA T(218), T(219), T(220), T(221), T(222), T(223), T(224),      WAN 1230
* T(225), T(226), T(227), T(228), T(229), T(230), T(231),      WAN 1240
* T(232) /.443317383947527D0,.432305826033741D0,                 WAN 1250
* .421229418017624D0,.410089821468717D0,.398888707435459D0,      WAN 1260
* .387627756194516D0,.376308656998716D0,.364933107823654D0,      WAN 1270
* .353502815112970D0,.342019493522372D0,.330484865662417D0,      WAN 1280
* .318900661840106D0,.307268619799319D0,.295590484460136D0,      WAN 1290
* .283868007657082D0/                                             WAN 1300
DATA T(233), T(234), T(235), T(236), T(237), T(238), T(239),      WAN 1310
* T(240), T(241), T(242), T(243), T(244), T(245), T(246),      WAN 1320
* T(247) /.272102947876337D0,.260297069991943D0,                 WAN 1330
* .248452145001057D0,.236569949758284D0,.224652266709132D0,      WAN 1340
* .212700883622626D0,.200717593323127D0,.188704193421389D0,      WAN 1350
* .176662486044902D0,.164594277567554D0,.152501378338656D0,      WAN 1360
* .140385602411376D0,.128248767270607D0,.116092693560333D0,      WAN 1370
* .103919204810509D0/                                             WAN 1380
DATA T(248), T(249), T(250), T(251), T(252), T(253), T(254),      WAN 1390
* T(255) /.917301271635196D-1,.795272891002330D-1,              WAN 1400
* .673125211657164D-1,.550876556946340D-1,.428545265363791D-1,      WAN 1410
* .306149687799790D-1,.183708184788137D-1,.612391237518953D-2/   WAN 1420
DATA W(1), W(2), W(3), W(4), W(5), W(6), W(7), W(8), W(9),      WAN 1430
* W(10), W(11), W(12), W(13), W(14), W(15)                        WAN 1440
* /1.0D0,.347854845137454D0,.652145154862546D0,                 WAN 1450
* .101228536290376D0,.222381034453374D0,.313706645877887D0,      WAN 1460
* .362683783378362D0,.271524594117541D-1,.622535239386479D-1,      WAN 1470
* .951585116824928D-1,.124628971255534D0,.149595988816577D0,      WAN 1480
* .169156519395033D0,.182603415044924D0,.189450610455068D0/     WAN 1490
DATA W(16), W(17), W(18), W(19), W(20), W(21), W(22), W(23),      WAN 1500
* W(24), W(25), W(26), W(27), W(28), W(29), W(30), W(31)         WAN 1510
* /.701861000947010D-2,.162743947309057D-1,.253920653092621D-1,      WAN 1520
* .342738629130214D-1,.428358980222267D-1,.509980592623762D-1,      WAN 1530
* .586840934785355D-1,.658222227763618D-1,.723457941088485D-1,      WAN 1540
* .781938957870703D-1,.833119242269468D-1,.876520930044038D-1,      WAN 1550
* .911738786957639D-1,.938443990808046D-1,.956387200792749D-1,      WAN 1560
* .965400885147278D-1/                                             WAN 1570
DATA W(32), W(33), W(34), W(35), W(36), W(37), W(38), W(39),      WAN 1580
* W(40), W(41), W(42), W(43), W(44), W(45), W(46), W(47)         WAN 1590
* /.178328072169643D-2,.414703326056247D-2,.650445796897836D-2,      WAN 1600
* .884675982636395D-2,.111681394601311D-1,.134630478967186D-1,      WAN 1610
* .157260304760247D-1,.179517157756973D-1,.201348231535302D-1,      WAN 1620
* .222701738083833D-1,.243527025687109D-1,.263774697150547D-1,      WAN 1630
* .283396726142595D-1,.302346570724025D-1,.320579283548516D-1,      WAN 1640
* .338051618371416D-1/                                             WAN 1650
DATA W(48), W(49), W(50), W(51), W(52), W(53), W(54), W(55),      WAN 1660
* W(56), W(57), W(58), W(59), W(60), W(61), W(62), W(63)         WAN 1670
* /.354722132568824D-1,.370551285402400D-1,.385501531786156D-1,      WAN 1680
* .399537411327203D-1,.412625632426235D-1,.424735151236536D-1,      WAN 1690
* .435837245293235D-1,.445905581637566D-1,.454916279274181D-1,      WAN 1700

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* .462847965813144D-1, .469681828162100D-1, .475401657148303D-1, WAN 1710
 * .479993885964583D-1, .483447622348030D-1, .485754674415034D-1, WAN 1720
 * .486909570091397D-1/ WAN 1730
 DATA W(64), W(65), W(66), W(67), W(68), W(69), W(70), W(71),
 * W(72), W(73), W(74), W(75), W(76), W(77), W(78), W(79) WAN 1740
 * /.449380960292090D-3, .104581267934035D-2, .164250301866903D-2, WAN 1750
 * .223828843096262D-2, .283275147145799D-2, .342552604091022D-2, WAN 1760
 * .401625498373864D-2, .460458425670296D-2, .519016183267633D-2, WAN 1770
 * .577263754286570D-2, .635166316170719D-2, .692689256689881D-2, WAN 1780
 * .749798192563473D-2, .806458989048606D-2, .862637779861675D-2, WAN 1790
 * .918300987166087D-2/ WAN 1800
 DATA W(80), W(81), W(82), W(83), W(84), W(85), W(86), W(87),
 * W(88), W(89), W(90), W(91), W(92), W(93), W(94), W(95) WAN 1810
 * /.973415341500681D-2, .102794790158322D-1, .108186607395031D-1, WAN 1820
 * .113513763240804D-1, .118773073727403D-1, .123961395439509D-1, WAN 1830
 * .129075627392673D-1, .134112712886163D-1, .139069641329520D-1, WAN 1840
 * .143943450041668D-1, .148731226021473D-1, .153430107688651D-1, WAN 1850
 * .158037286593993D-1, .162550009097852D-1, .166965578015892D-1, WAN 1860
 * .171281354231114D-1/ WAN 1870
 DATA W(96), W(97), W(98), W(99), W(100), W(101), W(102),
 * W(103), W(104), W(105), W(106), W(107), W(108), W(109), WAN 1880
 * W(110), W(111), W(112) /.175494758271177D-1, WAN 1890
 * .179603271850087D-1, .183604439373313D-1, .187495869405447D-1, WAN 1900
 * .191275236099509D-1, .194940280587066D-1, .198488812328309D-1, WAN 1910
 * .201918710421300D-1, .205227924869601D-1, .208414477807511D-1, WAN 1920
 * .211476464682213D-1, .214412055392085D-1, .217219495380521D-1, WAN 1930
 * .219897106684605D-1, .222443288937998D-1, .224856520327450D-1, WAN 1940
 * .227135358502365D-1/ WAN 1950
 DATA W(113), W(114), W(115), W(116), W(117), W(118), W(119),
 * W(120), W(121), W(122), W(123), W(124), W(125), W(126), WAN 1960
 * W(127) /.229278441436868D-1, .231284488243870D-1, WAN 1970
 * .233152299940628D-1, .234880760165359D-1, .236468835844476D-1, WAN 1980
 * .237915577810034D-1, .239220121367035D-1, .240381686810241D-1, WAN 1990
 * .241399579890193D-1, .242273192228152D-1, .243002001679719D-1, WAN 2000
 * .243585572646906D-1, .244023556338496D-1, .244315690978500D-1, WAN 2010
 * .244461801962625D-1/ WAN 2020
 DATA W(128), W(129), W(130), W(131), W(132), W(133), W(134),
 * W(135), W(136), W(137), W(138), W(139), W(140), W(141), WAN 2030
 * W(142) /.112789017822272D-3, .262534944296446D-3, WAN 2040
 * .412463254426176D-3, .562348954031410D-3, .712154163473321D-3, WAN 2050
 * .861853701420089D-3, .101142439320844D-2, .116084355756772D-2, WAN 2060
 * .131008868190250D-2, .145913733331073D-2, .160796713074933D-2, WAN 2070
 * .17565573633073D-2, .190488085349972D-2, .205292022796614D-2, WAN 2080
 * .220065164983991D-2/ WAN 2090
 DATA W(143), W(144), W(145), W(146), W(147), W(148), W(149),
 * W(150), W(151), W(152), W(153), W(154), W(155), W(156), WAN 2100
 * W(157) /.234805295632731D-2, .249510203470371D-2, WAN 2110
 * .264177682542749D-2, .278805532532771D-2, .293391559082972D-2, WAN 2120
 * .307933574119934D-2, .322429396179420D-2, .336876850731555D-2, WAN 2130
 * .351273770505631D-2, .365617995814250D-2, .379907374876626D-2, WAN 2140
 * .394139764140883D-2, .408313028605267D-2, .422425042138154D-2, WAN 2150
 * .436473687796806D-2/ WAN 2160
 DATA W(158), W(159), W(160), W(161), W(162), W(163), W(164),
 * W(165), W(166), W(167), W(168), W(169), W(170), W(171), WAN 2170
 * W(172) /.450456858144790D-2, .464372455568006D-2, WAN 2180
 * .478218392589269D-2, .491992592181387D-2, .505692988078684D-2, WAN 2190
 * .519317525086928D-2, .532864159391593D-2, .546330858864431D-2, WAN 2200
 * .559715603368291D-2, .573016385060144D-2, .586231208692265D-2, WAN 2210
 * .599358091911534D-2, .612395065556793D-2, .625340173954240D-2, WAN 2220
 * .638191475210788D-2/ WAN 2230
 DATA W(173), W(174), W(175), W(176), W(177), W(178), W(179),
 * W(180), W(181), W(182), W(183), W(184), W(185), W(186), WAN 2240
 * W(187) /.650947041505366D-2, .663604959378107D-2, WAN 2250
 * .676163330017380D-2, .688620269544632D-2, .700973909296982D-2, WAN 2260
 * .713222396107539D-2, .725363892583391D-2, .737396577381235D-2, WAN 2270
 * .749318645480588D-2, .761128308454566D-2, .772823794738156D-2, WAN 2280
 * .784403349893971D-2, .795865236875435D-2, .807207736287350D-2, WAN 2290
 * .818429146643827D-2/ WAN 2300
 DATA W(188), W(189), W(190), W(191), W(192), W(193), W(194),
 * W(195), W(196), W(197), W(198), W(199), W(200), W(201), WAN 2310
 * W(202) /.829527784623523D-2, .840501985322154D-2, WAN 2320
 * .85135002502249D-2, .862070508840101D-2, .872661596169881D-2, WAN 2330
 * .883121775724875D-2, .893449478375821D-2, .903643154866287D-2, WAN 2340
 * .913701276045081D-2, .92362233095630D-2, .933404837762327D-2, WAN 2350
 * .943047322573775D-2, .952548341062928D-2, .961906467984073D-2, WAN 2360
 * .971120299526628D-2/ WAN 2370
 * .981120299526628D-2/ WAN 2380
 * .991120299526628D-2/ WAN 2390
 * .101120299526628D-2/ WAN 2400
 * .102120299526628D-2/ WAN 2410
 * .103120299526628D-2/ WAN 2420
 * .104120299526628D-2/ WAN 2430
 * .105120299526628D-2/ WAN 2440
 * .106120299526628D-2/ WAN 2450
 * .107120299526628D-2/ WAN 2460

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DATA W(203), W(204), W(205), W(206), W(207), W(208), W(209),
* W(210), W(211), W(212), W(213), W(214), W(215), W(216),
* W(217) /.980188453525733D-2,.989109569669583D-2,
* .997882309703491D-2,.100650535763064D-1,.101497741990949D-1,
* .102329722564782D-1,.103146352679340D-1,.103947509832117D-1,
* .104733073841704D-1,.105502926865815D-1,.106256953418966D-1,
* .106995040389798D-1,.107717077058046D-1,.108422955111148D-1,
* .109112568660490D-1/
DATA W(218), W(219), W(220), W(221), W(222), W(223), W(224),
* W(225), W(226), W(227), W(228), W(229), W(230), W(231),
* W(232) /.109785814257296D-1,.110442590908139D-1,
* .111082800090098D-1,.111706345765534D-1,.112313134396497D-1,
* .112903074958755D-1,.113476078955455D-1,.114032060430392D-1,
* .114570935980906D-1,.115092624770395D-1,.115597048540436D-1,
* .116084131622531D-1,.116553800949452D-1,.117005986066207D-1,
* .117440619140606D-1/
DATA W(233), W(234), W(235), W(236), W(237), W(238), W(239),
* W(240), W(241), W(242), W(243), W(244), W(245), W(246),
* W(247) /.117857634973434D-1,.118256971008240D-1,
* .118638567340711D-1,.119002366727665D-1,.119348314595636D-1,
* .119676359049059D-1,.119986450878058D-1,.120278543565826D-1,
* .120552593295601D-1,.120808558957245D-1,.121046402153405D-1,
* .121266087205273D-1,.121467581157945D-1,.121650853785355D-1,
* .121815877594818D-1/
DATA W(248), W(249), W(250), W(251), W(252), W(253), W(254),
* W(255) /.121962627831147D-1,.122091082480372D-1,
* .122201222273040D-1,.122293030687103D-1,.122366493950402D-1,
* .122421601042728D-1,.122458343697479D-1,.122476716402898D-1/
LOOP = MAX0(1,N/256)
FLOOP = LOOP
H = (B-A)/FLOOP
SCALE = H/2.0D0
M = MIN0(128,N/2)
MT = 2*M
NPLACE = M - 1
DO 20 L=1,LOOP
  FL = L
  AL = A + (FL-1.0D0)*H
  BL = A + FL*H
  K = 256*(L-1)
  DO 10 I=1,M
    NPI = NPLACE + I
    S = T(NPI)
    R = W(NPI)*SCALE
    I1 = K + I
    I2 = K + MT + 1 - I
    TV(I1) = (AL*(1.0D0+S)+(1.0D0-S)*BL)/2.0D0
    TV(I2) = (AL*(1.0D0-S)+(1.0D0+S)*BL)/2.0D0
    WV(I1) = R
    WV(I2) = R
  10 CONTINUE
20 CONTINUE
RETURN
END

```

```

DOUBLE PRECISION FUNCTION RMIN(N, M, COND, UNITRD, AVERR)
C FOR A LINEAR SYSTEM (I-KMM)*XM=RHFCN OF ORDER M, THIS IS THE
C VALUE OF RELMIN USED IN IEQS. THE VARIABLE UNITRD IS DEFINED IN
C IEQAUS, AND THE VARIABLES COND AND AVERR ARE DEFINED IN IEQS
C USING CONEW.
C IT IS UNLIKELY THAT A SOLUTION X CAN BE FOUND FOR THE ORIGINAL
C INTEGRAL EQUATION WITH A SMALLER RELATIVE ERROR THAN RMIN.
DOUBLE PRECISION FLOAT1, FLOAT2, COND, AVERR, UNITRD, DMAX1
FLOAT1 = M
FLOAT2 = M/N
RMIN = DMAX1((FLOAT1**1.5D0)*COND*UNITRD, (FLOAT2**1.5D0)*
* AVERR)
RETURN
END

```

```

DOUBLE PRECISION FUNCTION CONEW(COND, ELINSY, RELRSD, AVERR,
* PASTC, PASTRE)
CON 10
* PASTC, PASTRE)
CON 20

```

```

C THIS IS USED IN UPDATING THE VALUE OF THE CONDITION NUMBER          CON 30
C IN IEQS.                                                            CON 40
  DOUBLE PRECISION COND, ELINSY, RELRSD, AVERR, PASTC, PASTRE,        CON 50
  * C, DSQRT, DMAX1                                                  CON 60
  AVERR = DSQRT(ELINSY*PASTRE)                                       CON 70
  PASTRE = ELINSY                                                    CON 80
  IF (RELRSD.EQ.0.0D0) GO TO 10                                       CON 90
  C = DMAX1(1.0D0,ELINSY/RELRSD)                                     CON 100
  CONEW = DSQRT(C*PASTC)                                             CON 110
  PASTC = C                                                         CON 120
  RETURN                                                            CON 130
10 CONEW = COND                                                    CON 140
  RETURN                                                            CON 150
  END                                                                CON 160

  SUBROUTINE LEAVE(IERSET, NF, MF, XV, TV, WV, ERROR, KERNEL,        LEA 10
  * RHFCN, EP, IFLAG, X, T, NT, IER, EPS, ELINSY, TN, WN, TM,      LEA 20
  * WM, XM, XMZ, KMM, KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN,    LEA 30
  * NUP, NHALF, XNORM)                                             LEA 40
C THIS ROUTINE SETS ALL NECESSARY PARAMETERS FOR LEAVING IEQAUS.   LEA 50
C IF NT .GT. 0, IT ALSO PERFORMS THE NECESSARY NYSTROM            LEA 60
C INTERPOLATION AT THE NODES GIVEN IN T.                          LEA 70
  DOUBLE PRECISION KERNEL, RHFCN, EP, X, T, XV, TV, WV, ERROR,    LEA 80
  * EPS, ELINSY, TN, WN, TM, WM, XM, XMZ, KMM, KMN, KNM, RHS,    LEA 90
  * IMKNN, LUFAC, R, RH, DELN, NORMK, R1, R2, FINLEP, SAVEP,      LEA 100
  * SUM, DERROR, NUMR1, NORM, TEMP, DABS, DMAX1, XNORM            LEA 110
  DIMENSION X(1), T(1), XV(MF), TV(MF), WV(MF), TN(NF), WN(NF),  LEA 120
  * WM(MF), XM(MF), XMZ(MF), KMM(NUP,NUP), KMN(NUP,NHALF),      LEA 130
  * KNM(NHALF,NUP), RHS(MF), IMKNN(NUP,NUP), LUFAC(NUP,NUP),    LEA 140
  * R(MF), RH(NF), TM(MF), DELN(NF)                               LEA 150
  COMMON /INFO/ R1, R2, FINLEP, NORMK, NFINAL, MFINAL            LEA 160
  EXTERNAL KERNEL, RHFCN                                         LEA 170
C SET ERROR PARAMETERS FOR RETURN.                                  LEA 180
  NORMK = 0.0D0                                                  LEA 190
  NFINAL = NF                                                    LEA 200
  MFINAL = MF                                                    LEA 210
  FINLEP = EPS                                                  LEA 220
  IF ((EPS.GT.EP) .AND. (ERROR.LE.EPS)) GO TO 10                 LEA 230
  IER = IERSET                                                  LEA 240
  EP = ERROR                                                    LEA 250
  IF (NT.EQ.0) GO TO 20                                         LEA 260
  GO TO 40                                                       LEA 270
10 IER = 3                                                       LEA 280
C SINCE EPS IS THE SMALLEST ERROR POSSIBLE, SET EP=EPS FOR THE    LEA 290
C RETURN ERROR ESTIMATE.                                          LEA 300
  EP = EPS                                                       LEA 310
  IF (NT.GT.0) GO TO 40                                         LEA 320
C NO NYSTROM INTERPOLATION IS DESIRED. RETURN THE VALUES AT THE  LEA 330
C GAUSSIAN NODE POINTS.                                          LEA 340
20 DO 30 I=1,MF                                                  LEA 350
  X(I) = XV(I)                                                  LEA 360
  T(I) = TV(I)                                                  LEA 370
30 CONTINUE                                                    LEA 380
  NT = MF                                                       LEA 390
  RETURN                                                        LEA 400
C CALCULATE NORM(K).                                             LEA 410
40 SAVEP = EP                                                  LEA 420
  DO 50 I=1,NF                                                  LEA 430
  IMKNN(I,I) = IMKNN(I,I) - 1.0D0                               LEA 440
50 CONTINUE                                                    LEA 450
  NORMK = 0.0D0                                                 LEA 460
  DO 70 I=1,NF                                                  LEA 470
  SUM = 0.0D0                                                  LEA 480
  DO 60 J=1,NF                                                  LEA 490
  SUM = SUM + DABS(IMKNN(I,J))                                  LEA 500
60 CONTINUE                                                    LEA 510
  NORMK = DMAX1(NORMK,SUM)                                       LEA 520
70 CONTINUE                                                    LEA 530
  DO 80 I=1,NF                                                  LEA 540
  IMKNN(I,I) = IMKNN(I,I) + 1.0D0                               LEA 550
80 CONTINUE                                                    LEA 560
  IF (NF.EQ.MF) GO TO 130                                       LEA 570
C ITERATE TO DECREASE THE NOISE LEVEL IN X. THIS SHOULD REDUCE   LEA 580
C POSSIBLE ERRORS IN NYSTROM INTERPOLATION.                     LEA 590

```



```

DERROR = ((1.0D0-R1)/R1)*EPS/NORMK
IF (IFLAG.EQ.1) DERROR = DERROR*XNORM
ITLOOP = 0
DO 90 I=1,MF
  XM(I) = XV(I)
90 CONTINUE
100 DO 110 I=1,MF
  XMZ(I) = XM(I)
110 CONTINUE
  CALL ITERT(KERNEL, RHFCN, NF, TN, WN, MF, TM, WM, XM, XMZ,
  * KMM, KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, NUP, NHALF,
  * 1)
  NUMR1 = NORM(XM, XMZ, MF, 1)
  ITLOOP = ITLOOP + 1
  IF ((NUMR1.GT.DERROR) .AND. (ITLOOP.LT.5)) GO TO 100
  DO 120 I=1,MF
    XV(I) = XM(I)
120 CONTINUE
C ESTIMATE NEW ERROR BOUND FOR NYSTROM INTERPOLATES.
  TEMP = NORMK*(R1/(1.0D0-R1))*NUMR1
  IF (IFLAG.EQ.1) TEMP = TEMP/XNORM
  EP = DMAX1(EP, TEMP)
  GO TO 140
C NO ITERATION USED IN COMPUTING X. JUST COMPUTE ERROR ESTIMATE IN
C NYSTROM INTERPOLATE.
130 TEMP = NORMK*ELINSY
  IF (IFLAG.EQ.0) TEMP = TEMP*XNORM
  IF (IER.NE.2) EP = DMAX1(EP, TEMP)
C COMPUTE NYSTROM INTERPOLATES AT THE NODES IN T.
140 DO 160 I=1,NT
  SUM = 0.0D0
  DO 150 J=1,MF
    SUM = SUM + WV(J)*KERNEL(T(I),TV(J))*XV(J)
150 CONTINUE
  X(I) = RHFCN(T(I)) + SUM
160 CONTINUE
  IF ((IER.EQ.0) .AND. (EP.GT.EPS)) IER = 4
  IF ((IER.EQ.1) .AND. (EP.GT.ERROR)) IER = 5
  IF ((IER.EQ.3) .AND. (EP.GT.EPS)) IER = 6
  EP = SAVEP
  RETURN
END

```

LEA 600
LEA 610
LEA 620
LEA 630
LEA 640
LEA 650
LEA 660
LEA 670
LEA 680
LEA 690
LEA 700
LEA 710
LEA 720
LEA 730
LEA 740
LEA 750
LEA 760
LEA 770
LEA 780
LEA 790
LEA 800
LEA 810
LEA 820
LEA 830
LEA 840
LEA 850
LEA 860
LEA 870
LEA 880
LEA 890
LEA 900
LEA 910
LEA 920
LEA 930
LEA 940
LEA 950
LEA 960
LEA 970
LEA 980
LEA 990
LEA 1000
LEA 1010

```

SUBROUTINE INTERP(TM, WM, XM, M, TN, WN, XN, N, KERNEL,
  * RHFCN, RHS, KMN, NHALF, NUP)
C USE THE VALUES OF XN(I), I=1,...,N, TO CALCULATE THE NYSTROM
C INTERPOLATES XM(I), I=1,...,M.
  DOUBLE PRECISION KERNEL, RHFCN, TM, WM, XM, TN, WN, XN, RHS,
  * KMN
  DIMENSION TM(M), WM(M), XM(M), TN(N), WN(N), XN(N), RHS(M),
  * KMN(NUP,NHALF)
  IF (M.GT.NUP) GO TO 60
C SINCE M .LE. NUPPER, SAVE K(TM(I),TN(J))=KMN(I,J) AND
C RHS(I)=RHFCN(TM(I)) FOR LATER USE IN ITERT.
  DO 20 I=1,M
    DO 10 J=1,N
      KMN(I,J) = WN(J)*KERNEL(TM(I),TN(J))
10 CONTINUE
20 CONTINUE
  DO 30 I=1,M
    RHS(I) = RHFCN(TM(I))
    XM(I) = RHS(I)
30 CONTINUE
C CALCULATE NYSTROM INTERPOLATING FORMULA.
  DO 50 I=1,M
    DO 40 J=1,N
      XM(I) = XM(I) + KMN(I,J)*XN(J)
40 CONTINUE
50 CONTINUE
  RETURN
C M .GT. NUPPER, SO SAVE JUST RHS(I) FOR LATER USE IN ITERT.
C CALCULATE NYSTROM INTERPOLATING FORMULA.
60 DO 80 I=1,M
  RHS(I) = RHFCN(TM(I))

```

INT 10
INT 20
INT 30
INT 40
INT 50
INT 60
INT 70
INT 80
INT 90
INT 100
INT 110
INT 120
INT 130
INT 140
INT 150
INT 160
INT 170
INT 180
INT 190
INT 200
INT 210
INT 220
INT 230
INT 240
INT 250
INT 260
INT 270
INT 280
INT 290
INT 300
INT 310

```

      XM(I) = RHS(I)
      DO 70 J=1,N
          XM(I) = XM(I) + WN(J)*KERNEL(TM(I),TN(J))*XN(J)
70  CONTINUE
80  CONTINUE
      RETURN
      END

```

SUBROUTINE ITERT(KERNEL, RHFCN, N, TN, WN, M, TM, WM, XM,
 * XMZ, KMM, KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, NUP,
 * NHALF, IFLG)
 C THIS ROUTINE CALCULATES ONE ITERATE XM GIVEN THE INITIAL GUESS
 C XMZ. THE ROUTINE IS DIVIDED ACCORDING TO WHETHER OR NOT
 C M .GT. NUPPER.
 DOUBLE PRECISION KERNEL, RHFCN, TN, WN, TM, WM, XM, XMZ, KMM,
 * KMN, KNM, RHS, IMKNN, LUFAC, R, RH, DELN, SUM, DET
 DIMENSION TN(N), WN(N), TM(M), WM(M), XM(M), XMZ(M),
 * KMM(NUP,NUP), KMN(NUP,NHALF), KNM(NHALF,NUP), RHS(M),
 * IMKNN(NUP,NUP), LUFAC(NUP,NUP), R(M), RH(N), DELN(N)
 C M .GT. NUPPER MEANS THAT THE MATRICES KMM,KMN,KNM CAN NO LONGER
 C BE STORED DUE TO LACK OF SPACE.
 IF (M.GT.NUP) GO TO 120
 IF (IFLG.EQ.1) GO TO 40
 C IF IFLG=0, THEN THE MATRICES KMM AND KNM MUST BE COMPUTED
 C AND STORED.
 DO 30 J=1,M
 DO 10 I=1,M
 KMM(I,J) = WM(J)*KERNEL(TM(I),TM(J))
10 CONTINUE
 DO 20 I=1,N
 KNM(I,J) = WM(J)*KERNEL(TN(I),TM(J))
20 CONTINUE
30 CONTINUE
 C COMPUTE RESIDUALS R(I)=RHFCN(TM(I))-XMZ(I)+KMM(TM(I))*XMZ(I)
 DO 40 I=1,M
 SUM = 0.0D0
 DO 50 J=1,M
 SUM = SUM + KMM(I,J)*XMZ(J)
50 CONTINUE
 R(I) = RHS(I) - (XMZ(I)-SUM)
60 CONTINUE
 C COMPUTE RH=KM*R AT ALL TN(I).
 DO 80 I=1,N
 RH(I) = 0.0D0
 DO 70 J=1,M
 RH(I) = RH(I) + KNM(I,J)*R(J)
70 CONTINUE
80 CONTINUE
 C CALCULATE DELN=((I-KN)**(-1))*KM*R AT ALL TN(I).
 C *****
 C
 CALL LINSYS(IMKNN, LUFAC, N, RH, DELN, 4, DET, NUP)
 C
 C SEE THE ORIGINAL REFERENCE IN IEQS.
 C *****
 C CALCULATE NEW XM.
 DO 110 I=1,M
 SUM = 0.0D0
 DO 90 J=1,M
 SUM = SUM + KMM(I,J)*R(J)
90 CONTINUE
 DO 100 J=1,N
 SUM = SUM + KMN(I,J)*DELN(J)
100 CONTINUE
 XM(I) = SUM + R(I) + XMZ(I)
110 CONTINUE
 RETURN
 C ENTRANCE WHEN M .GT. NUP.
 C CALCULATE RESIDUALS.
 DO 120 I=1,M
 SUM = 0.0D0
 DO 130 J=1,M
 SUM = SUM + WM(J)*KERNEL(TM(I),TM(J))*XMZ(J)
130 CONTINUE

```

      R(I) = RHS(I) - (XMZ(I)-SUM)
140 CONTINUE
C CALCULATE RH=KM*R.
      DO 160 I=1,N
          RH(I) = 0.0D0
          DO 150 J=1,M
              RH(I) = RH(I) + WM(J)*KERNEL(TN(I),TM(J))*R(J)
150 CONTINUE
160 CONTINUE
C *****
C
      CALL LINSYS(IMKNN, LUFAC, N, RH, DELN, 4, DET, NUP)
C
C SEE THE ORIGINAL REFERENCE IN IEQS.
C *****
C CALCULATE XM.
      DO 190 I=1,M
          SUM = 0.0D0
          DO 170 J=1,M
              SUM = SUM + WM(J)*KERNEL(TM(I),TM(J))*R(J)
170 CONTINUE
          DO 180 J=1,N
              SUM = SUM + WN(J)*KERNEL(TM(I),TN(J))*DELN(J)
180 CONTINUE
          XM(I) = SUM + R(I) + XMZ(I)
190 CONTINUE
      RETURN
      END
      ITE 670
      ITE 680
      ITE 690
      ITE 700
      ITE 710
      ITE 720
      ITE 730
      ITE 740
      ITE 750
      ITE 760
      *
      ITE 770
      ITE 780
      *
      ITE 790
      *
      ITE 800
      *
      ITE 810
      ITE 820
      ITE 830
      ITE 840
      ITE 850
      ITE 860
      ITE 870
      ITE 880
      ITE 890
      ITE 900
      ITE 910
      ITE 920
      ITE 930
      ITE 940

      DOUBLE PRECISION FUNCTION NORM(X, Y, N, IFLAG)
C IFLAG=0 CALCULATE THE MAXIMUM NORM OF X.
C IFLAG=1 CALCULATE THE MAXIMUM NORM OF X-Y.
      DOUBLE PRECISION X, Y, DMAX1, DABS
      DIMENSION X(N), Y(N)
      IF (IFLAG.EQ.1) GO TO 20
C FIND THE NORM OF X.
      NORM = 0.0D0
      DO 10 I=1,N
          NORM = DMAX1(NORM,DABS(X(I)))
10 CONTINUE
      RETURN
C FIND THE NORM OF X-Y.
20 NORM = 0.0D0
      DO 30 I=1,N
          NORM = DMAX1(NORM,DABS(X(I)-Y(I)))
30 CONTINUE
      RETURN
      END
      NOR 10
      NOR 20
      NOR 30
      NOR 40
      NOR 50
      NOR 60
      NOR 70
      NOR 80
      NOR 90
      NOR 100
      NOR 110
      NOR 120
      NOR 130
      NOR 140
      NOR 150
      NOR 160
      NOR 170
      NOR 180
      NOR 190

      SUBROUTINE LINSYS(A, D, N, B, X, OPTION, DET, MACHIN)
C SOLVE A*X=B, ORDER(A)=N, DIMENSION OF A=MACHIN.
C OPTION=1 SOLVE A*X=B, LEAVE THE LU DECOMPOSITION A=L*U IN D
C AND THE PIVOTS IN PIVOT. THE ANSWERS ARE LEFT IN X.
C IT IS PERMISSABLE TO LET B=X AND D=A, BUT THEN THE
C ORIGINAL CONTENTS OF A AND B ARE LOST.
C OPTION=2 CALCULATE DECOMPOSITION A=L*U INCLUDING PIVOTS. SOLVE
C A*X=B, AND THEN CALCULATE THE RESIDUAL AND ONE
C CORRECTION. THE CORRECTIONS ARE LEFT IN R, THE NEW
C VALUE OF X1 IN X, THE RELATIVE ERROR
C NORM(X0-X1)/NORM(X1)
C IN THE VARIABLE ERROR, AND THE RELATIVE RESIDUAL
C NORM(RESIDUAL)/NORM(B)
C IN THE VARIABLE RELRSD. THESE VALUES CAN BE OBTAINED
C USING THE COMMON/LINEAR/ GIVEN BELOW.
C OPTION=3 SAME AS OPTION=1, EXCEPT A=L*U IS KNOWN AND STORED
C IN D.
C OPTION=4 SAME AS OPTION=2, EXCEPT A=L*U IS KNOWN AND STORED
C IN D.
C THE DECOMPOSITION OF A INTO L*U USES SCALED PARTIAL PIVOTING IN
C THE COLUMNS. FOR OPTIONS 1 AND 2, THE DETERMINANT OF A IS
C CALCULATED AND STORED IN DET. IF DET=0, THEN THE ANSWERS ARE
C NONSENSE, D AND X DO NOT CONTAIN USEFUL INFORMATION, AND AND A
      LIN 10
      LIN 20
      LIN 30
      LIN 40
      LIN 50
      LIN 60
      LIN 70
      LIN 80
      LIN 90
      LIN 100
      LIN 110
      LIN 120
      LIN 130
      LIN 140
      LIN 150
      LIN 160
      LIN 170
      LIN 180
      LIN 190
      LIN 200
      LIN 210
      LIN 220
      LIN 230

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C AND B ARE LEFT UNDISTURBED (UNLESS D=A OR X=B).          LIN 240
C LINSYS IS PRESENTLY LIMITED TO N .LE. 100. TO REMOVE THIS LIN 250
C RESTRICTION,CHANGE THE DIMENSION STATEMENT FOR THE ARRAYS R, LIN 260
C SCALE, AND PIVOT, WHICH ARE GIVEN IN COMMON/LINEAR/.    LIN 270
  INTEGER OPTION, PIVOT                                  LIN 280
  DOUBLE PRECISION NORMX, NORME, NORMB, NORMR, A, D, B, X, R, LIN 290
  * SCALE, ERROR, RELRSD, DET, C, TEMP, DMAX1, DABS, SUM   LIN 300
  DIMENSION A(MACHIN,MACHIN), D(MACHIN,MACHIN), B(N), X(N) LIN 310
  COMMON /LINEAR/ R(100), SCALE(100), ERROR, RELRSD, PIVOT(100) LIN 320
  ISWIT = 1                                              LIN 330
  IF (OPTION.GT.2) GO TO 110                             LIN 340
C PRODUCE LU DECOMPOSITION OF A AND DET(A)                LIN 350
  DET = 0.0D0                                           LIN 360
  DO 20 I=1,N                                           LIN 370
    DO 10 J=1,N                                         LIN 380
      D(I,J) = A(I,J)                                  LIN 390
    10 CONTINUE                                        LIN 400
  20 CONTINUE                                           LIN 410
C PRODUCE SCALING FACTORS FOR SCALED PARTIAL PIVOTING.   LIN 420
  DO 40 I=1,N                                           LIN 430
    SCALE(I) = 0.0D0                                    LIN 440
    DO 30 J=1,N                                         LIN 450
      SCALE(I) = DMAX1(SCALE(I),DABS(D(I,J)))           LIN 460
    30 CONTINUE                                        LIN 470
    IF (SCALE(I).EQ.0.0D0) RETURN                       LIN 480
  40 CONTINUE                                           LIN 490
  DET = 1.0D0                                           LIN 500
  NM1 = N - 1                                           LIN 510
C BEGIN GAUSSIAN ELIMINATION.                             LIN 520
  DO 100 K=1,NM1                                        LIN 530
C SELECT PIVOT ROW.                                       LIN 540
  C = DABS(D(K,K))/SCALE(K)                             LIN 550
  INDEX = K                                             LIN 560
  KP1 = K + 1                                           LIN 570
  DO 50 I=KP1,N                                         LIN 580
    TEMP = DABS(D(I,K))/SCALE(I)                       LIN 590
    IF (TEMP.LE.C) GO TO 50                             LIN 600
    INDEX = I                                           LIN 610
    C = TEMP                                            LIN 620
  50 CONTINUE                                           LIN 630
  PIVOT(K) = INDEX                                     LIN 640
  IF (INDEX.EQ.K) GO TO 70                              LIN 650
C SWITCH ROWS OF D.                                       LIN 660
  DO 60 J=K,N                                           LIN 670
    TEMP = D(K,J)                                       LIN 680
    D(K,J) = D(INDEX,J)                                LIN 690
    D(INDEX,J) = TEMP                                  LIN 700
  60 CONTINUE                                           LIN 710
  TEMP = SCALE(K)                                       LIN 720
  SCALE(K) = SCALE(INDEX)                              LIN 730
  SCALE(INDEX) = TEMP                                  LIN 740
  DET = -DET                                           LIN 750
  70 DET = DET*D(K,K)                                   LIN 760
C ELIMINATE UNKNOWN =K FROM BELOW DIAGONAL IN COLUMN K.  LIN 770
  IF (DET.EQ.0.0D0) RETURN                             LIN 780
  DO 90 I=KP1,N                                         LIN 790
    D(I,K) = D(I,K)/D(K,K)                             LIN 800
    TEMP = D(I,K)                                       LIN 810
    DO 80 J=KP1,N                                       LIN 820
      D(I,J) = D(I,J) - TEMP*D(K,J)                   LIN 830
    80 CONTINUE                                        LIN 840
  90 CONTINUE                                           LIN 850
  100 CONTINUE                                          LIN 860
  DET = DET*D(N,N)                                       LIN 870
  IF (DET.EQ.0.0D0) RETURN                             LIN 880
C THE DECOMPOSITION A=P*L*U IS COMPLETE.                 LIN 890
C BEGIN SOLUTION OF LINEAR SYSTEM P*L*U*X=B             LIN 900
C SET R AND X FOR SOLUTION OF A*X=B.                   LIN 910
  110 DO 120 I=1,N                                       LIN 920
    R(I) = B(I)                                         LIN 930
    X(I) = 0.0D0                                        LIN 940
  120 CONTINUE                                          LIN 950
  GO TO 170                                             LIN 960
C SET R=B-A*X FOR COMPUTATION OF CORRECTION.           LIN 970
  130 DO 150 I=1,N                                       LIN 980
    SUM = 0.0D0                                        LIN 990

```

```

DO 140 J=1,N
    SUM = SUM + A(I,J)*X(J)
140 CONTINUE
    R(I) = B(I) - SUM
150 CONTINUE
    NORMB = 0.0D0
    NORMR = 0.0D0
    DO 160 I=1,N
        NORMB = DMAX1(NORMB,DABS(B(I)))
        NORMR = DMAX1(NORMR,DABS(R(I)))
160 CONTINUE
    RELRSD = NORMR/NORMB
    ISWIT = 2
C SOLVE P*L*Z=R AND STORE IN R.
170 NM1 = N - 1
    DO 200 K=1,NM1
        IF (PIVOT(K).EQ.K) GO TO 180
        KPIV = PIVOT(K)
        TEMP = R(K)
        R(K) = R(KPIV)
        R(KPIV) = TEMP
180 KPI = K + 1
        DO 190 I=KPI,N
            R(I) = R(I) - D(I,K)*R(K)
190 CONTINUE
200 CONTINUE
C SOLVE U*E=R AND STORE IN R. ALSO LET X=X+E.
    R(N) = R(N)/D(N,N)
    GO TO (210, 220, 210, 220), OPTION
210 X(N) = R(N)
    GO TO 230
220 X(N) = X(N) + R(N)
230 DO 270 II=1,NM1
    I = N - II
    SUM = 0.0D0
    IP1 = I + 1
    DO 240 J=IP1,N
        SUM = SUM + D(I,J)*R(J)
240 CONTINUE
    R(I) = (R(I)-SUM)/D(I,I)
    GO TO (250, 260, 250, 260), OPTION
250 X(I) = R(I)
    GO TO 270
260 X(I) = X(I) + R(I)
270 CONTINUE
C SOLUTION OF LINEAR SYSTEM IS COMPLETE. RETURN FOR OPTION=1,3.
    GO TO (280, 290, 280, 290), OPTION
280 RETURN
290 GO TO (130, 300), ISWIT
C CALCULATE ERRORS BASED ON CORRECTION.
300 NORMX = 0.0D0
    NORME = 0.0D0
    DO 310 I=1,N
        NORMX = DMAX1(NORMX,DABS(X(I)))
        NORME = DMAX1(NORME,DABS(R(I)))
310 CONTINUE
    ERROR = NORME/NORMX
    RETURN
    END

```

```

LIN 1000
LIN 1010
LIN 1020
LIN 1030
LIN 1040
LIN 1050
LIN 1060
LIN 1070
LIN 1080
LIN 1090
LIN 1100
LIN 1110
LIN 1120
LIN 1130
LIN 1140
LIN 1150
LIN 1160
LIN 1170
LIN 1180
LIN 1190
LIN 1200
LIN 1210
LIN 1220
LIN 1230
LIN 1240
LIN 1250
LIN 1260
LIN 1270
LIN 1280
LIN 1290
LIN 1300
LIN 1310
LIN 1320
LIN 1330
LIN 1340
LIN 1350
LIN 1360
LIN 1370
LIN 1380
LIN 1390
LIN 1400
LIN 1410
LIN 1420
LIN 1430
LIN 1440
LIN 1450
LIN 1460
LIN 1470
LIN 1480
LIN 1490
LIN 1500
LIN 1510
LIN 1520
LIN 1530
LIN 1540
LIN 1550
LIN 1560
LIN 1570
LIN 1580

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ALGORITHM 504

GERK: Global Error Estimation for Ordinary Differential Equations [D]

L. F. SHAMPINE and H. A. WATTS

Sandia Laboratories

Key Words and Phrases: ordinary differential equations, initial value problems, global error estimation, Runge-Kutta-Fehlberg method, Fortran code GERK

CR Categories: 3.20, 5.17

Language: Fortran

DESCRIPTION

This algorithm is a complement to [1] where the theoretical development is described.

REFERENCES

- SHAMPINE, L.F., AND WATTS, H.A. Global error estimation for ordinary differential equations. *ACM Trans. Math. Software* 2, 2 (June 1976), 172-186.

ALGORITHM

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SUBROUTINE GERK(F, NEQN, Y, T, TOUT, RELERR, ABSERR, IFLAG,      GER  10
* GERROR, WORK, IWORK)                                       GER  20
C                                                             GER  30
C   FEHLBERG FOURTH (FIFTH) ORDER RUNGE-KUTTA METHOD WITH    GER  40
C   GLOBAL ERROR ASSESSMENT                                  GER  50
C                                                             GER  60
C   WRITTEN BY H.A.WATTS AND L.F.SHAMPINE                    GER  70
C   SANDIA LABORATORIES                                     GER  80
C                                                             GER  90
C   GERK IS DESIGNED TO SOLVE SYSTEMS OF DIFFERENTIAL EQUATIONS GER 100
C   WHEN IT IS IMPORTANT TO HAVE A READILY AVAILABLE GLOBAL ERROR GER 110
C   ESTIMATE. PARALLEL INTEGRATION IS PERFORMED TO YIELD TWO GER 120
C   SOLUTIONS ON DIFFERENT MESH SPACINGS AND GLOBAL EXTRAPOLATION GER 130
C   IS APPLIED TO PROVIDE AN ESTIMATE OF THE GLOBAL ERROR IN THE GER 140
C   MORE ACCURATE SOLUTION.                                  GER 150
C                                                             GER 160
C   FOR IBM SYSTEM 360 AND 370 AND OTHER MACHINES OF SIMILAR GER 170
C   ARITHMETIC CHARACTERISTICS, THIS CODE SHOULD BE CONVERTED TO GER 180
C   DOUBLE PRECISION.                                       GER 190
C                                                             GER 200
C*****                                                       GER 210
C ABSTRACT                                                    GER 220
C*****                                                       GER 230
C   SUBROUTINE GERK INTEGRATES A SYSTEM OF NEQN FIRST ORDER GER 240
C   ORDINARY DIFFERENTIAL EQUATIONS OF THE FORM              GER 250
C   DY(I)/DT = F(T,Y(1),Y(2),...,Y(NEQN))                   GER 260
C   WHERE THE Y(I) ARE GIVEN AT T.                           GER 270
C   TYPICALLY THE SUBROUTINE IS USED TO INTEGRATE FROM T TO TOUT GER 280
C   BUT IT CAN BE USED AS A ONE-STEP INTEGRATOR TO ADVANCE THE GER 290
C   SOLUTION A SINGLE STEP IN THE DIRECTION OF TOUT. ON RETURN, AN GER 300
C   GER 310

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ACM Transactions on Mathematical Software, Vol. 2, No. 2, June 1976, Pages 200-203.

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C ESTIMATE OF THE GLOBAL ERROR IN THE SOLUTION AT T IS PROVIDED GER 320
C AND THE PARAMETERS IN THE CALL LIST ARE SET FOR CONTINUING THE GER 330
C INTEGRATION. THE USER HAS ONLY TO CALL GERK AGAIN (AND PERHAPS GER 340
C DEFINE A NEW VALUE FOR TOUT). ACTUALLY, GERK IS MERELY AN GER 350
C INTERFACING ROUTINE WHICH ALLOCATES VIRTUAL STORAGE IN THE GER 360
C ARRAYS WORK, IWORK AND CALLS SUBROUTINE GERKS FOR THE SOLUTION. GER 370
C GERKS IN TURN CALLS SUBROUTINE FEHL WHICH COMPUTES AN APPROX- GER 380
C IMATE SOLUTION OVER ONE STEP. GER 390
C GER 400
C GERK USES THE RUNGE-KUTTA-FEHLBERG (4,5) METHOD DESCRIBED GER 410
C IN THE REFERENCE GER 420
C E.FEHLBERG , LOW-ORDER CLASSICAL RUNGE-KUTTA FORMULAS WITH GER 430
C STEPSIZE CONTROL , NASA TR K-315 GER 440
C GER 450
C GER 460
C THE PARAMETERS REPRESENT- GER 470
C F -- SUBROUTINE F(T,Y,YP) TO EVALUATE DERIVATIVES GER 480
C YP(I)=DY(I)/DT GER 490
C NEQN -- NUMBER OF EQUATIONS TO BE INTEGRATED GER 500
C Y(*) -- SOLUTION VECTOR AT T GER 510
C T -- INDEPENDENT VARIABLE GER 520
C TOUT -- OUTPUT POINT AT WHICH SOLUTION IS DESIRED GER 530
C RELERR,ABSERR -- RELATIVE AND ABSOLUTE ERROR TOLERANCES FOR GER 540
C LOCAL ERROR TEST. AT EACH STEP THE CODE REQUIRES THAT GER 550
C ABS(LOCAL ERROR) .LE. RELERR*ABS(Y) + ABSERR GER 560
C FOR EACH COMPONENT OF THE LOCAL ERROR AND SOLUTION GER 570
C VECTORS. GER 580
C IFLAG -- INDICATOR FOR STATUS OF INTEGRATION GER 590
C GERROR(*) -- VECTOR WHICH ESTIMATES THE GLOBAL ERROR AT T. GER 600
C THAT IS, GERROR(I) APPROXIMATES Y(I)-TRUE GER 610
C SOLUTION(I). GER 620
C WORK(*) -- ARRAY TO HOLD INFORMATION INTERNAL TO GERK WHICH GER 630
C IS NECESSARY FOR SUBSEQUENT CALLS. MUST BE DIMENSIONED GER 640
C AT LEAST 3+8*NEQN GER 650
C IWORK(*) -- INTEGER ARRAY USED TO HOLD INFORMATION INTERNAL GER 660
C TO GERK WHICH IS NECESSARY FOR SUBSEQUENT CALLS. MUST GER 670
C BE DIMENSIONED AT LEAST 5 GER 680
C GER 690
C GER 700
C***** GER 710
C FIRST CALL TO GERK GER 720
C***** GER 730
C GER 740
C THE USER MUST PROVIDE STORAGE IN HIS CALLING PROGRAM FOR THE GER 750
C ARRAYS IN THE CALL LIST - Y(NEQN), WORK(3+8*NEQN), IWORK(5), GER 760
C DECLARE F IN AN EXTERNAL STATEMENT, SUPPLY SUBROUTINE F(T,Y,YP) GER 770
C AND INITIALIZE THE FOLLOWING PARAMETERS- GER 780
C GER 790
C NEQN -- NUMBER OF EQUATIONS TO BE INTEGRATED. (NEQN .GE. 1) GER 800
C Y(*) -- VECTOR OF INITIAL CONDITIONS GER 810
C T -- STARTING POINT OF INTEGRATION , MUST BE A VARIABLE GER 820
C TOUT -- OUTPUT POINT AT WHICH SOLUTION IS DESIRED. GER 830
C T=TOUT IS ALLOWED ON THE FIRST CALL ONLY, IN WHICH CASE GER 840
C GERK RETURNS WITH IFLAG=2 IF CONTINUATION IS POSSIBLE. GER 850
C RELERR,ABSERR -- RELATIVE AND ABSOLUTE LOCAL ERROR TOLERANCES GER 860
C WHICH MUST BE NON-NEGATIVE BUT MAY BE CONSTANTS. WE CAN GER 870
C USUALLY EXPECT THE GLOBAL ERRORS TO BE SOMEWHAT SMALLER GER 880
C THAN THE REQUESTED LOCAL ERROR TOLERANCES. TO AVOID GER 890
C LIMITING PRECISION DIFFICULTIES THE CODE ALWAYS USES GER 900
C THE LARGER OF RELERR AND AN INTERNAL RELATIVE ERROR GER 910
C PARAMETER WHICH IS MACHINE DEPENDENT. GER 920
C IFLAG -- +1,-1 INDICATOR TO INITIALIZE THE CODE FOR EACH NEW GER 930
C PROBLEM. NORMAL INPUT IS +1. THE USER SHOULD SET IFLAG= GER 940
C -1 ONLY WHEN ONE-STEP INTEGRATOR CONTROL IS ESSENTIAL. GER 950
C IN THIS CASE, GERK ATTEMPTS TO ADVANCE THE SOLUTION A GER 960
C SINGLE STEP IN THE DIRECTION OF TOUT EACH TIME IT IS GER 970
C CALLED. SINCE THIS MODE OF OPERATION RESULTS IN EXTRA GER 980
C COMPUTING OVERHEAD, IT SHOULD BE AVOIDED UNLESS NEEDED. GER 990
C GER 1000
C GER 1010
C***** GER 1020
C OUTPUT FROM GERK GER 1030
C***** GER 1040
C GER 1050
C Y(*) -- SOLUTION AT T GER 1060
C T -- LAST POINT REACHED IN INTEGRATION. GER 1070
C IFLAG = 2 -- INTEGRATION REACHED TOUT. INDICATES SUCCESSFUL GER 1080
C RETURN AND IS THE NORMAL MODE FOR CONTINUING GER 1090
C INTEGRATION. GER 1100
C =-2 -- A SINGLE SUCCESSFUL STEP IN THE DIRECTION OF GER 1110
C TOUT HAS BEEN TAKEN. NORMAL MODE FOR CONTINUING GER 1120
C INTEGRATION ONE STEP AT A TIME. GER 1130
C = 3 -- INTEGRATION WAS NOT COMPLETED BECAUSE MORE THAN GER 1140
C 9000 DERIVATIVE EVALUATIONS WERE NEEDED. THIS GER 1150
C IS APPROXIMATELY 500 STEPS. GER 1160
C = 4 -- INTEGRATION WAS NOT COMPLETED BECAUSE SOLUTION GER 1170
C VANISHED MAKING A PURE RELATIVE ERROR TEST GER 1180
C IMPOSSIBLE. MUST USE NON-ZERO ABSERR TO CONTINUE. GER 1190
C USING THE ONE-STEP INTEGRATION MODE FOR ONE STEP GER 1200
C IS A GOOD WAY TO PROCEED. GER 1210

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C          = 5 -- INTEGRATION WAS NOT COMPLETED BECAUSE REQUESTED          GER 1220
C          ACCURACY COULD NOT BE ACHIEVED USING SMALLEST                   GER 1230
C          ALLOWABLE STEPSIZE. USER MUST INCREASE THE ERROR                GER 1240
C          TOLERANCE BEFORE CONTINUED INTEGRATION CAN BE                   GER 1250
C          ATTEMPTED.                                                       GER 1260
C          = 6 -- GERK IS BEING USED INEFFICIENTLY IN SOLVING              GER 1270
C          THIS PROBLEM. TOO MUCH OUTPUT IS RESTRICTING THE                GER 1280
C          NATURAL STEPSIZE CHOICE. USE THE ONE-STEP                       GER 1290
C          INTEGRATOR MODE.                                                 GER 1300
C          = 7 -- INVALID INPUT PARAMETERS                                  GER 1310
C          THIS INDICATOR OCCURS IF ANY OF THE FOLLOWING IS                 GER 1320
C          SATISFIED - NEQN .LE. 0                                          GER 1330
C          T=TOUT AND IFLAG .NE. +1 OR -1                                   GER 1340
C          RELERR OR ABSERR .LT. 0.                                         GER 1350
C          IFLAG .EQ. 0 OR .LT. -2 OR .GT. 7                                GER 1360
C          GERROR(*) -- ESTIMATE OF THE GLOBAL ERROR IN THE SOLUTION AT T    GER 1370
C          WORK(*), IWORK(*) -- INFORMATION WHICH IS USUALLY OF NO          GER 1380
C          INTEREST TO THE USER BUT NECESSARY FOR SUBSEQUENT              GER 1390
C          CALLS. WORK(1),... ,WORK(NEQN) CONTAIN THE FIRST                 GER 1400
C          DERIVATIVES OF THE SOLUTION VECTOR Y AT T.                     GER 1410
C          WORK(NEQN+1) CONTAINS THE STEPSIZE H TO BE                      GER 1420
C          ATTEMPTED ON THE NEXT STEP. IWORK(1) CONTAINS                   GER 1430
C          THE DERIVATIVE EVALUATION COUNTER.                               GER 1440
C                                                                           GER 1450
C                                                                           GER 1460
C*****                                                                    GER 1470
C SUBSEQUENT CALLS TO GERK                                                GER 1480
C*****                                                                    GER 1490
C                                                                           GER 1500
C SUBROUTINE GERK RETURNS WITH ALL INFORMATION NEEDED TO CONTINUE         GER 1510
C THE INTEGRATION. IF THE INTEGRATION REACHED TOUT, THE USER NEED        GER 1520
C ONLY DEFINE A NEW TOUT AND CALL GERK AGAIN. IN THE ONE-STEP             GER 1530
C INTEGRATOR MODE (IFLAG=-2) THE USER MUST KEEP IN MIND THAT EACH        GER 1540
C STEP TAKEN IS IN THE DIRECTION OF THE CURRENT TOUT. UPON               GER 1550
C REACHING TOUT (INDICATED BY CHANGING IFLAG TO 2), THE USER MUST       GER 1560
C THEN DEFINE A NEW TOUT AND RESET IFLAG TO -2 TO CONTINUE IN THE        GER 1570
C ONE-STEP INTEGRATOR MODE.                                              GER 1580
C                                                                           GER 1590
C IF THE INTEGRATION WAS NOT COMPLETED BUT THE USER STILL WANTS         GER 1600
C TO CONTINUE (IFLAG=3 CASE), HE JUST CALLS GERK AGAIN. THE              GER 1610
C FUNCTION COUNTER IS THEN RESET TO 0 AND ANOTHER 9000 FUNCTION          GER 1620
C EVALUATIONS ARE ALLOWED.                                               GER 1630
C                                                                           GER 1640
C HOWEVER, IN THE CASE IFLAG=4, THE USER MUST FIRST ALTER THE           GER 1650
C ERROR CRITERION TO USE A POSITIVE VALUE OF ABSERR BEFORE              GER 1660
C INTEGRATION CAN PROCEED. IF HE DOES NOT, EXECUTION IS TERMINATED.     GER 1670
C                                                                           GER 1680
C ALSO, IN THE CASE IFLAG=5, IT IS NECESSARY FOR THE USER TO            GER 1690
C RESET IFLAG TO 2 (OR -2 WHEN THE ONE-STEP INTEGRATION MODE IS          GER 1700
C BEING USED) AS WELL AS INCREASING EITHER ABSERR, RELERR OR BOTH        GER 1710
C BEFORE THE INTEGRATION CAN BE CONTINUED. IF THIS IS NOT DONE,         GER 1720
C EXECUTION WILL BE TERMINATED. THE OCCURRENCE OF IFLAG=5              GER 1730
C INDICATES A TROUBLE SPOT (SOLUTION IS CHANGING RAPIDLY,              GER 1740
C SINGULARITY MAY BE PRESENT) AND IT OFTEN IS INADVISABLE TO           GER 1750
C CONTINUE.                                                                GER 1760
C                                                                           GER 1770
C IF IFLAG=6 IS ENCOUNTERED, THE USER SHOULD USE THE ONE-STEP          GER 1780
C INTEGRATION MODE WITH THE STEPSIZE DETERMINED BY THE CODE. IF          GER 1790
C THE USER INSISTS UPON CONTINUING THE INTEGRATION WITH GERK IN         GER 1800
C THE INTERVAL MODE, HE MUST RESET IFLAG TO 2 BEFORE CALLING GERK       GER 1810
C AGAIN. OTHERWISE, EXECUTION WILL BE TERMINATED.                        GER 1820
C                                                                           GER 1830
C IF IFLAG=7 IS OBTAINED, INTEGRATION CAN NOT BE CONTINUED UNLESS       GER 1840
C THE INVALID INPUT PARAMETERS ARE CORRECTED.                             GER 1850
C                                                                           GER 1860
C IT SHOULD BE NOTED THAT THE ARRAYS WORK, IWORK CONTAIN                GER 1870
C INFORMATION REQUIRED FOR SUBSEQUENT INTEGRATION. ACCORDINGLY,           GER 1880
C WORK AND IWORK SHOULD NOT BE ALTERED.                                  GER 1890
C                                                                           GER 1900
C*****                                                                    GER 1910
C                                                                           GER 1920
C          DIMENSION Y(NEQN), GERROR(NEQN), WORK(1), IWORK(5)             GER 1930
C          EXTERNAL F                                                       GER 1940
C COMPUTE INDICES FOR THE SPLITTING OF THE WORK ARRAY                    GER 1950
C          K1M = NEQN + 1                                                    GER 1960
C          K1 = K1M + 1                                                       GER 1970
C          K2 = K1 + NEQN                                                     GER 1980
C          K3 = K2 + NEQN                                                     GER 1990
C          K4 = K3 + NEQN                                                     GER 2000
C          K5 = K4 + NEQN                                                     GER 2010
C          K6 = K5 + NEQN                                                     GER 2020
C          K7 = K6 + NEQN                                                     GER 2030
C          K8 = K7 + NEQN                                                     GER 2040
C*****                                                                    GER 2050
C THIS INTERFACING ROUTINE MERELY RELIEVES THE USER OF A LONG           GER 2060
C CALLING LIST VIA THE SPLITTING APART OF TWO WORKING STORAGE           GER 2070
C ARRAYS. IF THIS IS NOT COMPATIBLE WITH THE USERS COMPILER,           GER 2080
C HE MUST USE GERKS DIRECTLY.                                             GER 2090
C*****                                                                    GER 2100
C          CALL GERKS(F, NEQN, Y, T, TOUT, RELERR, ABSERR, IFLAG,         GER 2110
C          * GERROR, WORK(1), WORK(K1M), WORK(K1), WORK(K2), WORK(K3),     GER 2120

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* WORK(K4), WORK(K5), WORK(K6), WORK(K7), WORK(K8),          GER 2130
* WORK(K8+1), IWORK(1), IWCRK(2), IWORK(3), IWORK(4), IWORK(5) GER 2140
RETURN                                                       GER 2150
END                                                           GER 2160

SUBROUTINE GERKS(F, NEQN, Y, T, TOUT, RELERR, ABSEK, IFLAG,      GER 10
* GERROR, YP, H, F1, F2, F3, F4, F5, YG, YGP, SAVRE, SAVAE,    GER 20
* NFE, KOP, INIT, JFLAG, KFLAG)                               GER 30
C FEHLBERG FOURTH(FIFTH) ORDER RUNGE-KUTTA METHOD WITH        GER 40
C GLOBAL ERKOR ASSESSMENT                                     GER 50
C *****                                                     GER 60
C GERKS INTEGRATES A SYSTEM OF FIRST ORDER ORDINARY DIFFERENTIAL GER 70
C EQUATIONS AS DESCRIBED IN THE COMMENTS FOR GERK. THE ARRAYS GER 80
C YP,F1,F2,F3,F4,F5,YG AND YGP (OF DIMENSION AT LEAST NEQN) AND GER 90
C THE VARIABLES H,SAVRE,SAVAE,NFE,KOP,INIT,JFLAG,AND KFLAG ARE GER 100
C USED INTERNALLY BY THE CODE AND APPEAR IN THE CALL LIST TO GER 110
C ELIMINATE LOCAL RETENTION OF VARIABLES BETWEEN CALLS.      GER 120
C ACCORDINGLY, THEY SHOULD NOT BE ALTERED. ITEMS OF POSSIBLE GER 130
C INTEREST ARE                                               GER 140
C YP - DERIVATIVE OF SOLUTION VECTOR AT T                     GER 150
C H - AN APPROPRIATE STEPSIZE TO BE USED FOR THE NEXT STEP   GER 160
C NFE- COUNTER ON THE NUMBER OF DERIVATIVE FUNCTION           GER 170
C EVALUATIONS.                                               GER 180
C *****                                                     GER 190
C LOGICAL HFAILD, OUTPUT                                     GER 200
C DIMENSION Y(NEQN), YP(NEQN), F1(NEQN), F2(NEQN), F3(NEQN), GER 210
C * F4(NEQN), F5(NEQN), YG(NEQN), YGP(NEQN), GERROR(NEQN)    GER 220
C EXTERNAL F                                                 GER 230
C *****                                                     GER 240
C THE COMPUTER UNIT ROUND OFF ERROR U IS THE SMALLEST POSITIVE VALUE GER 250
C REPRESENTABLE IN THE MACHINE SUCH THAT 1.+ U .GT. 1.        GER 260
C VALUES TO BE USED ARE                                     GER 270
C U = 9.5E-7 FOR IBM 360/370                                  GER 280
C U = 1.5E-8 FOR UNIVAC 1108                                  GER 290
C U = 7.5E-9 FOR PDP-10                                       GER 300
C U = 7.1E-15 FOR CDC 6000 SERIES                              GER 310
C U = 2.2E-16 FOR IBM 360/370 DOUBLE PRECISION                GER 320
C DATA U / /                                               GER 330
C *****                                                     GER 340
C REMIN IS A TOLERANCE THRESHOLD WHICH IS ALSO DETERMINED BY THE GER 350
C INTEGRATION METHOD. IN PARTICULAR, A FIFTH ORDER METHOD WILL GER 360
C GENERALLY NOT BE CAPABLE OF DELIVERING ACCURACIES NEAR LIMITING GER 370
C PRECISION ON COMPUTERS WITH LONG WORDLENGTHS.              GER 380
C DATA REMIN /3.E-11/                                       GER 390
C *****                                                     GER 400
C THE EXPENSE IS CONTROLLED BY RESTRICTING THE NUMBER        GER 410
C OF FUNCTION EVALUATIONS TO BE APPROXIMATELY MAXNFE.        GER 420
C AS SET, THIS CORRESPONDS TO ABOUT 500 STEPS.               GER 430
C DATA MAXNFE /9000/                                         GER 440
C *****                                                     GER 450
C CHECK INPUT PARAMETERS                                     GER 460
C IF (NEQN.LT.1) GO TO 10                                       GER 470
C IF ((RELERR.LT.0.) .OR. (ABSEK.LT.0.)) GO TO 10             GER 480
C MFLAG = IAES(IFLAG)                                          GER 490
C IF ((MFLAG.GE.1) .AND. (MFLAG.LE.7)) GO TO 20              GER 500
C INVALID INPUT                                             GER 510
C 10 IFLAG = 7                                               GER 520
C RETURN                                                     GER 530
C IS THIS THE FIRST CALL                                     GER 540
C 20 IF (MFLAG.EQ.1) GO TO 70                                  GER 550
C CHECK CONTINUATION POSSIBILITIES                           GER 560
C IF (T.EQ.TOUT) GO TO 10                                     GER 570
C IF (MFLAG.NE.2) GO TO 30                                    GER 580
C IFLAG = +2 OR -2                                           GER 590
C IF (INIT.EQ.0) GO TO 60                                       GER 600
C IF (KFLAG.EQ.3) GO TO 50                                       GER 610
C IF ((KFLAG.EQ.4) .AND. (ABSERR.EQ.0.)) GO TO 40            GER 620
C IF ((KFLAG.EQ.5) .AND. (RELERR.LE.SAVRE) .AND.            GER 630
* (ABSERR.LE.SAVAE)) GO TO 40                                     GER 640
C GO TO 70                                                    GER 650
C IFLAG = 3,4,5,6, OR 7                                       GER 660
C 30 IF (IFLAG.EQ.3) GO TO 50                                       GER 670
C IF ((IFLAG.EQ.4) .AND. (ABSEK.GT.0.)) GO TO 60            GER 680
C INTEGRATION CANNOT BE CONTINUED SINCE USER DID NOT RESPOND TO GER 690
C THE INSTRUCTIONS PERTAINING TO IFLAG=4,5,6 OR 7           GER 700
C 40 STOP                                                     GER 710
C *****                                                     GER 720
C RESET FUNCTION EVALUATION COUNTER                           GER 730
C 50 NFE = 0                                                 GER 740
C IF (MFLAG.EQ.2) GO TO 70                                       GER 750
C RESET FLAG VALUE FROM PREVIOUS CALL                         GER 760
C 60 IFLAG = JFLAG                                           GER 770
C SAVE INPUT IFLAG AND SET CONTINUATION FLAG VALUE FOR SUBSEQUENT GER 780
C INPUT CHECKING                                             GER 790
C 70 JFLAG = IFLAG                                           GER 800
C KFLAG = 0                                                  GER 810
C SAVE RELERR AND ABSERR FOR CHECKING INPUT ON SUBSEQUENT CALLS GER 820
C SAVRE = RELERR                                             GER 830
C SAVAE = ABSEK                                             GER 840

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C RESTRICT RELATIVE ERROR TOLERANCE TO BE AT LEAST AS LARGE AS      GER 850
C 32U+REMIN TO AVOID LIMITING PRECISION DIFFICULTIES ARISING        GER 860
C FROM IMPOSSIBLE ACCURACY REQUESTS                                  GER 870
      RER = AMAX1(RELERR,32.*U+REMIN)                                  GER 880
      U26 = 26.*U                                                    GER 890
      DT = TOUT - T                                                  GER 900
      IF (MFLAG.EQ.1) GO TO 80                                        GER 910
      IF (INIT.EQ.0) GO TO 90                                        GER 920
      GO TO 110                                                       GER 930
C *****                                                             GER 940
C      INITIALIZATION --                                           GER 950
C      SET INITIALIZATION COMPLETION INDICATOR,INIT                 GER 960
C      SET INDICATOR FOR TOO MANY OUTPUT POINTS,KOP                 GER 970
C      EVALUATE INITIAL DERIVATIVES                                  GER 980
C      COPY INITIAL VALUES AND DERIVATIVES FOR THE                 GER 990
C      PARALLL SOLUTION                                             GER 1000
C      SET COUNTER FOR FUNCTION EVALUATIONS,NFE                     GER 1010
C      ESTIMATE STARTING STEPSIZE                                    GER 1020
      80 INIT = 0                                                    GER 1030
      KOP = 0                                                         GER 1040
      A = T                                                            GER 1050
      CALL F(A, Y, YP)                                               GER 1060
      NFE = 1                                                         GER 1070
      IF (T.NE.TOUT) GO TO 90                                         GER 1080
      IFLAG = 2                                                       GER 1090
      RETURN                                                           GER 1100
      90 INIT = 1                                                    GER 1110
      H = ABS(DT)                                                     GER 1120
      TOLN = 0.                                                       GER 1130
      DO 100 K=1,NEQN                                                 GER 1140
        YG(K) = Y(K)                                                 GER 1150
        YGP(K) = YP(K)                                               GER 1160
        TOL = RER*ABS(Y(K)) + ABSERR                                  GER 1170
        IF (TOL.LE.0.) GO TO 100                                       GER 1180
        TOLN = TOL                                                   GER 1190
        YPK = ABS(YP(K))                                              GER 1200
        IF (YPK*H**5.GT.TOL) H = (TOL/YPK)**0.2                     GER 1210
      100 CONTINUE                                                    GER 1220
        IF (TOLN.LE.0.) H = 0.                                        GER 1230
        H = AMAX1(H,U26*AMAX1(ABS(T),ABS(DT)))                       GER 1240
C *****                                                             GER 1250
C      SET STEPSIZE FOR INTEGRATION IN THE DIRECTION FROM T TO TOUT GER 1260
      110 H = SIGN(H,DT)                                              GER 1270
C TEST TO SEE IF GERK IS BEING SEVERELY IMPACTED BY TOO MANY      GER 1280
C OUTPUT POINTS                                                    GER 1290
      IF (ABS(H).GT.2.*ABS(DT)) KOP = KOP + 1                       GER 1300
      IF (KOP.NE.100) GO TO 120                                       GER 1310
      KOP = 0                                                         GER 1320
      IFLAG = 6                                                       GER 1330
      RETURN                                                           GER 1340
      120 IF (ABS(DT).GT.U26*ABS(T)) GO TO 140                       GER 1350
C IF TOO CLOSE TO OUTPUT POINT,EXTRAPOLATE AND RETURN            GER 1360
      DO 130 K=1,NEQN                                                 GER 1370
        YG(K) = YG(K) + DT*YGP(K)                                     GER 1380
        Y(K) = Y(K) + DT*YP(K)                                       GER 1390
      130 CONTINUE                                                    GER 1400
      A = TOUT                                                         GER 1410
      CALL F(A, YG, YGP)                                             GER 1420
      CALL F(A, Y, YP)                                               GER 1430
      NFE = NFE + 2                                                  GER 1440
      GO TO 230                                                       GER 1450
C INITIALIZE OUTPUT POINT INDICATOR                                GER 1460
      140 OUTPUT = .FALSE.                                           GER 1470
C TO AVOID PREMATURE UNDERFLOW IN THE ERROR TOLERANCE FUNCTION,  GER 1480
C SCALE THE ERROR TOLERANCES                                       GER 1490
      SCALE = 2./RER                                                 GER 1500
      AE = SCALE*ABSERR                                              GER 1510
C *****                                                             GER 1520
C *****                                                             GER 1530
C      STEP BY STEP INTEGRATION                                     GER 1540
      150 HFAILD = .FALSE.                                           GER 1550
C SET SMALLEST ALLOWABLE STEPSIZE                                  GER 1560
      HMIN = U26*ABS(T)                                              GER 1570
C ADJUST STEPSIZE IF NECESSARY TO HIT THE OUTPUT POINT.          GER 1580
C LOOK AHEAD TWO STEPS TO AVOID DRASTIC CHANGES IN THE STEPSIZE GER 1590
C AND THUS LESSEN THE IMPACT OF OUTPUT POINTS ON THE CODE.       GER 1600
      DT = TOUT - T                                                  GER 1610
      IF (ABS(DT).GE.2.*ABS(H)) GO TO 170                             GER 1620
      IF (ABS(DT).GT.ABS(H)) GO TO 160                               GER 1630
C THE NEXT SUCCESSFUL STEP WILL COMPLETE THE INTEGRATION TO THE  GER 1640
C OUTPUT POINT                                                    GER 1650
      OUTPUT = .TRUE.                                               GER 1660
      H = DT                                                         GER 1670
      GO TO 170                                                       GER 1680
      160 H = 0.5*DT                                                 GER 1690
C *****                                                             GER 1700
C      CORE INTEGRATOR FOR TAKING A SINGLE STEP                    GER 1710
C *****                                                             GER 1720
C      THE TOLERANCES HAVE BEEN SCALED TO AVOID PREMATURE UNDERFLOW GER 1730
C      IN COMPUTING THE ERROR TOLERANCE FUNCTION ET.              GER 1740
C      TO AVOID PROBLEMS WITH ZERO CROSSINGS, RELATIVE ERROR IS  GER 1750

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C *****
C      SHOULD WE TAKE ANOTHER STEP
C      IF (OUTPUT) GO TO 230
C      IF (IFLAG.GT.0) GO TO 150
C *****
C      INTEGRATION SUCCESSFULLY COMPLETED
C      ONE-STEP MODE
C      IFLAG = -2
C      GO TO 240
C INTERVAL MODE
230 T = TOUT
      IFLAG = 2
240 DO 250 K=1,NEQN
      GERROR(K) = (YG(K)-Y(K))/31.
250 CONTINUE
      RETURN
      END
GER 2670
GER 2680
GER 2690
GER 2700
GER 2710
GER 2720
GER 2730
GER 2740
GER 2750
GER 2760
GER 2770
GER 2780
GER 2790
GER 2800
GER 2810
GER 2820
GER 2830
GER 2840

      SUBROUTINE FEHL(F, NEQN, Y, T, H, YP, F1, F2, F3, F4, F5, S)
C      FEHLBERG FOURTH-FIFTH ORDER RUNGE-KUTTA METHOD
C *****
C      FEHL INTEGRATES A SYSTEM OF NEQN FIRST ORDER
C      ORDINARY DIFFERENTIAL EQUATIONS OF THE FORM
C      DY(I)/DT=F(T,Y(1),---,Y(NEQN))
C      WHERE THE INITIAL VALUES Y(I) AND THE INITIAL DERIVATIVES
C      YP(I) ARE SPECIFIED AT THE STARTING POINT T. FEHL ADVANCES
C      THE SOLUTION OVER THE FIXED STEP H AND RETURNS
C      THE FIFTH ORDER (SIXTH ORDER ACCURATE LOCALLY) SOLUTION
C      APPROXIMATION AT T+H IN ARRAY S(I).
C      F1,---,F5 ARE ARRAYS OF DIMENSION NEQN WHICH ARE NEEDED
C      FOR INTERNAL STORAGE.
C      THE FORMULAS HAVE BEEN GROUPED TO CONTROL LOSS OF SIGNIFICANCE.
C      FEHL SHOULD BE CALLED WITH AN H NOT SMALLER THAN 13 UNITS OF
C      ROUND OFF IN T SO THAT THE VARIOUS INDEPENDENT ARGUMENTS CAN BE
C      DISTINGUISHED.
C *****
      DIMENSION Y(NEQN), YP(NEQN), F1(NEQN), F2(NEQN), F3(NEQN),
      * F4(NEQN), F5(NEQN), S(NEQN)
      CH = 0.25*H
      DO 10 K=1,NEQN
      F5(K) = Y(K) + CH*YP(K)
10 CONTINUE
      CALL F(T+0.25*H, F5, F1)
      CH = 0.09375*H
      DO 20 K=1,NEQN
      F5(K) = Y(K) + CH*(YP(K)+3.*F1(K))
20 CONTINUE
      CALL F(T+0.375*H, F5, F2)
      CH = H/2197.
      DO 30 K=1,NEQN
      F5(K) = Y(K) + CH*(1932.*YP(K)+(7296.*F2(K)-7200.*F1(K)))
30 CONTINUE
      CALL F(T+12./13.*H, F5, F3)
      CH = H/4104.
      DO 40 K=1,NEQN
      F5(K) = Y(K) + CH*((8341.*YP(K)-845.*F3(K))+(29440.*F2(K)
      * -32832.*F1(K)))
40 CONTINUE
      CALL F(T+H, F5, F4)
      CH = H/20520.
      DO 50 K=1,NEQN
      F1(K) = Y(K) + CH*((-6080.*YP(K)+(9295.*F3(K)-5643.*F4(K))
      * +(41040.*F1(K)-28352.*F2(K)))
50 CONTINUE
      CALL F(T+0.5*H, F1, F5)
C COMPUTE APPROXIMATE SOLUTION AT T+H
      CH = H/7618050.
      DO 60 K=1,NEQN
      S(K) = Y(K) + CH*((902860.*YP(K)+(3855735.*F3(K)-1371249.*
      * F4(K)))+(3953664.*F2(K)+277020.*F5(K)))
60 CONTINUE
      RETURN
      END
FEH 10
FEH 20
FEH 30
FEH 40
FEH 50
FEH 60
FEH 70
FEH 80
FEH 90
FEH 100
FEH 110
FEH 120
FEH 130
FEH 140
FEH 150
FEH 160
FEH 170
FEH 180
FEH 190
FEH 200
FEH 210
FEH 220
FEH 230
FEH 240
FEH 250
FEH 260
FEH 270
FEH 280
FEH 290
FEH 300
FEH 310
FEH 320
FEH 330
FEH 340
FEH 350
FEH 360
FEH 370
FEH 380
FEH 390
FEH 400
FEH 410
FEH 420
FEH 430
FEH 440
FEH 450
FEH 460
FEH 470
FEH 480
FEH 490
FEH 500
FEH 510
FEH 520
FEH 530
FEH 540
FEH 550

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ALGORITHM 505

A List Insertion Sort for Keys With Arbitrary Key Distribution [S20]

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Key Words and Phrases: sorting, searching, linked lists, lists, insertion

CR Categories: 5.31, 4.49, 3.74

Language: Basic Fortran

DESCRIPTION

The description of this algorithm, the experimental results, and the references are given in the author's paper, "A List Insertion Sort for Keys With Arbitrary Key Distribution," *ACM Trans. Math. Software* 2, 2(June 1976), 143-153.

ALGORITHM

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SUBROUTINE SPN(K, L, II, JJ, MIN)
C **SAMPLE SEARCH SORT**
C A LIST INSERTION SORT TO BUILD UP A SINGLE CIRCULARLY
C LINKED LIST.
C THIS ALGORITHM SORTS THE KEYS K(II),K(II+1),...,K(JJ) IN
C ASCENDING SORTING ORDER BY THE MEANS OF THE POINTER FIELDS
C L(II),L(II+1),...,L(JJ).
C THE EXPECTED NUMBER OF NECESSARY COMPARISONS AND
C ACCESSES IS OF ORDER N**1.5 (N=JJ-II+1).
C WITH SINGLE KEYS THE ALGORITHM IS PROPORTIONAL IN SORTING
C TIME TO A SHELL SORT ALGORITHM.
C ISTRT AUXILIARY VARIABLE IN RANDOM SEARCH
C K ARRAY OF N KEYS.
C L ARRAY FOR N POINTERS.
C IPOI VARIABLE WHICH IS USED TO MANIPULATE POINTERS.
C KDIFF DISTANCE OF THE KEY FOUND IN RANDOM SEARCH
C AND THE KEY WHICH SHALL BE INSERTED.
C KEY,MADIN ADDRESSES USED IN SEQUENTIAL AND RANDOM SEARCH.
C MAX,MIN VARIABLES TO STORE THE ADDRESS OF THE
C KEYS WITH THE SMALLEST AND LARGEST VALUE.
C N NUMBER OF KEYS.
C IRT STEP WIDTH IN RANDOM SEARCH.
C TAB ARRAY OF VALUES WHICH DETERMINES THE SAMPLE SIZE.
C THE ALGORITHM IS STABLE.
C MIN RETURNS THE ADDRESS OF THE FIRST RECORD IN SORTING
C ORDER OF THE LIST.
INTEGER TAB(57)
DIMENSION K(1), L(1)
DATA TAB(1), TAB(2), TAB(3), TAB(4), TAB(5), TAB(6), TAB(7),
* TAB(8), TAB(9), TAB(10), TAB(11), TAB(12), TAB(13), TAB(14),
* TAB(15), TAB(16), TAB(17), TAB(18), TAB(19), TAB(20),
* TAB(21), TAB(22), TAB(23), TAB(24), TAB(25), TAB(26),
* TAB(27), TAB(28), TAB(29), TAB(30), TAB(31), TAB(32),
* TAB(33), TAB(34), TAB(35), TAB(36), TAB(37), TAB(38),
* TAB(39), TAB(40), TAB(41), TAB(42), TAB(43), TAB(44),
* TAB(45), TAB(46), TAB(47), TAB(48), TAB(49), TAB(50),

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* TAB(51), TAB(52), TAB(53), TAB(54), TAB(55), TAB(56),
* TAB(57) /3,5,10,17,26,37,50,65,82,101,122,145,170,197,226,
* 257,290,325,362,401,442,485,530,577,626,677,730,785,842,901,
* 962,1025,1090,1157,1226,1297,1370,1445,1522,1601,1682,1765,
* 1850,1937,2026,2117,2210,2305,2402,2501,2602,2705,2810,2917,
* 3026,3137,3250/
C INITIALISATION (PART I)
C -----
MIN = II
MAX = II
L(II) = II
IF (JJ-II) 170, 170, 10
10 KMIN = K(MIN)
KMAX = KMIN
IA = II + 1
IRT = 1
ITAB = TAB(1) + II - 1
ISTRT = II
C -----
C INSERTION LOOP (PART II)
DO 160 J=IA,JJ
C CORRECTION OF THE SAMPLE SIZE (IF NECESSARY)
IF (J-ITAB) 30, 20, 20
C -----
20 IRT = IRT + 1
ISTRT = ISTRT + 1
ITAB = TAB(IRT) + II - 1
C -----
30 KJ = K(J)
C PROVISION FOR ARISING LARGEST AND SMALLEST KEYS (1).
IF (KJ-KMAX) 40, 60, 60
C -----
40 IF (KJ-KMIN) 70, 50, 90
C -----
C KEY WHICH SHALL BE INSERTED IS EQUAL TO THE MINIMAL
C KEY SORTED SO FAR (2).
C -----
50 MADIN = MIN
MIN = J
GO TO 130
C -----
C KEY WHICH SHALL BE INSERTED IS EQUAL TO OR LARGER THAN THE
C MAXIMAL KEY SORTED SO FAR (3).
C -----
60 I = MAX
MAX = J
KMAX = K(J)
GO TO 80
C -----
C KEY FOR INSERTION IS SMALLER THAN THE SMALLEST KEY SORTED
C SO FAR (4).
C -----
70 I = MAX
MIN = J
KMIN = K(J)
C -----
C INSERTION OF THE RECORD AND CORRECTION OF THE VALUE WHICH
C DETERMINES THE SAMPLE SIZE (5).
C -----
80 IPOI = L(I)
L(I) = J
L(J) = IPOI
GO TO 160
C -----
C INITIALISATION FOR RANDOM SEARCH (FIRST STEP) (6).
C -----
90 MADIN = MIN
IPOI = J - 1
I = KJ - KMIN
C -----
C RANDOM SEARCH (7).
C PART 1 (7A).
C -----
DO 120 KEY=ISTRT,IPOI,IRT
KDIFF = KJ - K(KEY)
IF (KDIFF) 120, 140, 100
C -----
100 IF (I-KDIFF) 120, 120, 110
C -----
C PART 2.. STORE NEAREST KEY FOUND SO FAR (7B)
C -----
110 I = KDIFF
MADIN = KEY
C -----
120 CONTINUE
C -----
C SEQUENTIAL SEARCH (8).
C -----
130 IPOI = MADIN
MADIN = L(IPOI)
IF (KJ-K(MADIN)) 150, 130, 130

```

```

SPN 370
SPN 380
SPN 390
SPN 400
SPN 410
SPN 420
SPN 430
SPN 440
SPN 450
SPN 460
SPN 470
SPN 480
SPN 490
SPN 500
SPN 510
SPN 520
SPN 530
SPN 540
SPN 550
SPN 560
SPN 570
SPN 580
SPN 590
SPN 600
SPN 610
SPN 620
SPN 630
SPN 640
SPN 650
SPN 660
SPN 670
SPN 680
SPN 690
SPN 700
SPN 710
SPN 720
SPN 730
SPN 740
SPN 750
SPN 760
SPN 770
SPN 780
SPN 790
SPN 800
SPN 810
SPN 820
SPN 830
SPN 840
SPN 850
SPN 860
SPN 870
SPN 880
SPN 890
SPN 900
SPN 910
SPN 920
SPN 930
SPN 940
SPN 950
SPN 960
SPN 970
SPN 980
SPN 990
SPN 1000
SPN 1010
SPN 1020
SPN 1030
SPN 1040
SPN 1050
SPN 1060
SPN 1070
SPN 1080
SPN 1090
SPN 1100
SPN 1110
SPN 1120
SPN 1130
SPN 1140
SPN 1150
SPN 1160
SPN 1170
SPN 1180
SPN 1190
SPN 1200
SPN 1210
SPN 1220
SPN 1230
SPN 1240
SPN 1250
SPN 1260
SPN 1270

```



```
C -----
C KEY SEARCHED FOR WAS FOUND IN RANDOM SEARCH
C -----
140 MADIN = KEY
      GO TO 130
C -----
C INSERT (9)
C -----
150 L(IPOI) = J
      L(J) = MADIN
C -----
160 CONTINUE
C -----
170 MIN = L(MAX)
      L(MAX) = 0
      RETURN
      END
```

```
SPN 1280
SPN 1290
SPN 1300
SPN 1310
SPN 1320
SPN 1330
SPN 1340
SPN 1350
SPN 1360
SPN 1370
SPN 1380
SPN 1390
SPN 1400
SPN 1410
SPN 1420
SPN 1430
SPN 1440
```


ALGORITHM 506

HQR3 and EXCHNG: Fortran Subroutines for Calculating and Ordering the Eigenvalues of a Real Upper Hessenberg Matrix [F2]

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Key Words and Phrases: eigenvalue, QR-algorithm
CR Categories: 5.14
Language: Fortran

DESCRIPTION

1. Usage

HQR3 is a Fortran subroutine to reduce a real upper Hessenberg matrix A to quasi-triangular form B by a unitary similarity transformation U :

$$B = U^T A U.$$

The diagonal of B consists of 1×1 and 2×2 blocks as illustrated below:

$$\begin{bmatrix} x & x & x & x & x & x \\ 0 & x & x & x & x & x \\ 0 & x & x & x & x & x \\ 0 & 0 & 0 & x & x & x \\ 0 & 0 & 0 & 0 & x & x \\ 0 & 0 & 0 & 0 & x & x \end{bmatrix}$$

The 1×1 blocks contain the real eigenvalues of A , and the 2×2 blocks contain the complex eigenvalues, a conjugate pair to each block. The blocks are ordered so that the eigenvalues appear in descending order of absolute value along the diagonal. The transformation U is postmultiplied into an array V , which presumably contains earlier transformations performed on A .

The decomposition produced by *HQR3* differs from the one produced by the *EISPACK* subroutine *HQR2* [2] in that the eigenvalues of the final quasi-triangular matrix are ordered. This ordering makes the decomposition essentially unique, which is important in some applications (e.g. see [4]). It should also be noted that when the eigenvalues $\lambda_1, \lambda_2, \dots, \lambda_n$ of A are ordered so that $|\lambda_1| \geq |\lambda_2| \geq \dots \geq |\lambda_n|$, and if $|\lambda_i| > |\lambda_{i+1}|$, then the first i columns of U form an orthonormal basis for the invariant subspace corresponding to $\lambda_1, \lambda_2, \dots, \lambda_i$. In applications where it is desired to work with the matrix A in such a dominant invariant subspace,

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HQR3 provides a convenient means for calculating a basis. The corresponding leading principal submatrix of B is a representation of A in that subspace. When an ordered quasi-triangular form is not required, the *EISPACK* program *HQR2* will be slightly more efficient than *HQR3*.

The calling sequence for *HQR3* is:

CALL HQR3 (A,V,N,NLOW,NUP,EPS,ER,EI,TYPE,NA,NV)

with (parameters preceded by an asterisk are altered by the subroutine):

- *A A doubly subscripted real array containing the matrix to be reduced. On return, A contains the final quasi-triangular matrix. The elements of the array below the third subdiagonal are unaltered by the subroutine.
- *V A doubly subscripted real array into which the reducing transformation is postmultiplied.
- N An integer containing the order of A and V.
- NLOW } Integers prescribing what part of A is to be reduced. Specifically
- NUP } A(NLOW,NLOW-1) and A(NUP+1,NUP) are assumed zero and only the block from NLOW through NUP is reduced. However, the transformation is performed on all of the matrix A so that the final result is similar to A.
- EPS Convergence criterion. Maximal accuracy will be attained if EPS is set to β^{-t} , where β is the base of the floating-point word and t is the length of its fraction. Smaller values of EPS will increase the amount of work without significantly improving the accuracy.
- *ER A singly subscripted real array containing the real parts of the eigenvalues.
- *EI A singly subscripted real array containing the imaginary parts of the eigenvalues.
- *TYPE A singly subscripted integer array whose i th entry is:
 - 0 if the i th eigenvalue is real;
 - 1 if the i th eigenvalue is complex with positive imaginary part;
 - 2 if the i th eigenvalue is complex with negative imaginary part;
 - 1 if the i th eigenvalue was not successfully calculated.
 The entry 1 is always followed by a 2. Only elements NLOW through NUP of ER, EI, and TYPE are set by *HQR3*.
- NA The first dimension of the array A.
- NV The first dimension of the array V.

HQR3 can be used together with the *EISPACK* programs *ORTHES* and *ORTRAN* [2] to reduce a full matrix A to quasi-triangular form with the eigenvalues appearing in descending order of absolute value along the diagonal,

```
CALL ORTHES(NA,N,NLOW,NUP,A,P)
CALL ORTRAN(NA,N,NLOW,NUP,A,P,V)
CALL HQR3(A,V,N,NLOW,NUP,EPS,ER,EI,TYPE,NA,NA)
```

where P is a singly subscripted scratch array of order N.

HQR3 requires the subroutines *EXCHNG*, *SPLIT*, and *QRSTEP*.

EXCHNG is a Fortran subroutine to interchange consecutive 1×1 and 2×2 blocks of an upper Hessenberg matrix. Specifically it is supposed that the upper Hessenberg matrix A has a block of order $b1$ starting at the l th diagonal element and a block of order $b2$ starting at the $(l + b1)$ -th diagonal element (illustrated below for $n = 5, l = 2, b1 = 2, b2 = 1$):

```

      l
      x x x x x
l 0 x x x x
  0 x x x x
  0 0 0 x x
  0 0 0 0 x
```

EXCHNG produces an orthogonal similarity transformation W such that $W^T A W$ has consecutive blocks of order $b2$ and $b1$ starting at the l th diagonal element (illustrated from the example above):

```

      l
      x x x x x
l 0 x x x x
  0 0 x x x
  0 0 x x x
  0 0 0 0 x
```

The eigenvalues associated with each block are interchanged along with the blocks. The transformation \bar{W} is postmultiplied into the matrix V .

EXCHNG can be used to rearrange the blocks of the quasi-triangular matrix produced by *HQR3*. For example, one might wish to cluster a group of nearly equal eigenvalues at the top of the matrix before applying a deflation technique to uncouple them from the rest of the problem.

The calling sequence for *EXCHNG* is:

```
CALL EXCHNG(A,V,N,L,B1,B2,EPS,FAIL,NB,NV)
```

with (parameters preceded by an asterisk are altered by the subroutine):

- *A A doubly subscripted real array containing the matrix whose blocks are to be interchanged. Only rows and columns L through $L + B1 + B2 - 1$ are transformed. Elements of the array A below the third subdiagonal are not altered.
- *V A doubly subscripted array containing the matrix V into which the reducing transformation is to be accumulated. Only columns L through $L + B1 + B2 - 1$ are altered.
- N An integer containing the order of A and V .
- L An integer containing the leading diagonal position of the blocks.
- B1 An integer containing the size of the first block.
- B2 An integer containing the size of the second block.
- EPS A convergence criterion (cf. EPS in the calling sequence for *HQR3*).
- *FAIL A logical variable that on normal return is false. If the iteration to interchange the blocks failed, it is set to true.
- NA The first dimension of the array A .
- NV The first dimension of the array V .

By repeated applications of *EXCHNG* the eigenvalues of a quasi-triangular matrix can be arranged in any desired order.

EXCHNG requires the subroutine *QRSTEP*.

2. Method and Programming Details

HQR3 uses the implicit double-shift *QR* algorithm to reduce A to quasi-triangular form (for the theoretical background see [3, 5]). The program is essentially a Fortran variant of part of the Algol program *hqr2* in the *Numerische Mathematik* handbook series [1] with these differences:

- (1) The calling sequence is somewhat different.
- (2) Eigenvectors are not computed.
- (3) The parameters *NLOW* and *NUP* are not comparable to the parameters *low* and *upp* of *hqr2*. Specifically, in *HQR3* rows 1 through N of V are transformed; whereas in *hqr2* only rows *low* through *upp* are transformed.
- (4) The code that performs the *QR* iterations in *hqr2* is replaced by a call to the subroutine *QRSTEP*.
- (5) After a 1×1 or 2×2 block has been isolated, the subroutine *EXCHNG* is used to position it correctly among the previously isolated blocks. It should be realized that *EXCHNG* is a numerical algorithm and may change an ill-conditioned eigenvalue in a block significantly. This means that after interchanging two blocks with ill-conditioned eigenvalues that are very nearly equal in absolute value, the eigenvalues may still not be in descending order of absolute value. Since numerically the absolute values of these eigenvalues cannot be told apart, this phenomenon is not of much importance.

The convergence criterion is the same for both programs. A subdiagonal element $A(I + 1, I)$ is regarded as negligible if it is less than or equal to $\text{EPS} * (\text{ABS}(A(I, I)) + \text{ABS}(A(I + 1, I + 1)))$. If 10 or 20 iterations are performed without convergence, an ad hoc shift is introduced to break up any cycling. If 30 iterations are performed without convergence, the subroutine gives up. Although when this happens the matrix returned is not quasi-triangular, it is still almost exactly similar to the original matrix, and the similarity transformation has been accumulated in V .

EXCHNG works as follows. The first block is used to determine an implicit *QR* shift. An arbitrary *QR* step is performed on both blocks to eliminate the uncoupling between them. Then a sequence of *QR* steps using the previously determined shift

is performed on both blocks. Except in ill-conditioned cases, a block of size $B1$ having the eigenvalues of the first block will emerge in the lower part of the array occupied by both blocks, usually in one or two steps. If 30 iterations pass without convergence (the criterion is the same as in *HQR3*), the subroutine gives an error return.

Both *HQR3* and *EXCHNG* use the subroutine *QRSTEP* to perform the *QR* iterations. In addition, *HQR3* uses the subroutine *SPLIT* to separate real eigenvalues of a 2×2 block.

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ALGORITHM

SUBROUTINE HQR3(A, V, N, NLOW, NUP, EPS, ER, EI, TYPE, NA, NV)	HQR	10
INTEGER N, NA, NLOW, NUP, NV, TYPE(N)	HQR	20
REAL A(NA,N), EI(N), ER(N), EPS, V(NV,N)	HQR	30
C HQR3 REDUCES THE UPPER HESSENBERG MATRIX A TO QUASI-	HQR	40
C TRIANGULAR FORM BY UNITARY SIMILARITY TRANSFORMATIONS.	HQR	50
C THE EIGENVALUES OF A, WHICH ARE CONTAINED IN THE 1X1	HQR	60
C AND 2X2 DIAGONAL BLOCKS OF THE REDUCED MATRIX, ARE	HQR	70
C ORDERED IN DESCENDING ORDER OF MAGNITUDE ALONG THE	HQR	80
C DIAGONAL. THE TRANSFORMATIONS ARE ACCUMULATED IN THE	HQR	90
C ARRAY V. HQR3 REQUIRES THE SUBROUTINES EXCHNG,	HQR	100
C QRSTEP, AND SPLIT. THE PARAMETERS IN THE CALLING	HQR	110
C SEQUENCE ARE (STARRED PARAMETERS ARE ALTERED BY THE	HQR	120
C SUBROUTINE)	HQR	130
C *A AN ARRAY THAT INITIALLY CONTAINS THE N X N	HQR	140
C UPPER HESSENBERG MATRIX TO BE REDUCED. ON	HQR	150
C RETURN A CONTAINS THE REDUCED, QUASI-	HQR	160
C TRIANGULAR MATRIX.	HQR	170
C *V AN ARRAY THAT CONTAINS A MATRIX INTO WHICH	HQR	180
C THE REDUCING TRANSFORMATIONS ARE TO BE	HQR	190
C MULTIPLIED.	HQR	200
C N THE ORDER OF THE MATRICES A AND V.	HQR	210
C NLOW A(NLOW,NLOW-1) AND A(NUP,+1,NUP) ARE	HQR	220
C NUP ASSUMED TO BE ZERO, AND ONLY ROWS NLOW	HQR	230
C THROUGH NUP AND COLUMNS NLOW THROUGH	HQR	240
C NUP ARE TRANSFORMED, RESULTING IN THE	HQR	250
C CALCULATION OF EIGENVALUES NLOW	HQR	260
C THROUGH NUP.	HQR	270
C EPS A CONVERGENCE CRITERION.	HQR	280
C *ER AN ARRAY THAT ON RETURN CONTAINS THE REAL	HQR	290
C PARTS OF THE EIGENVALUES.	HQR	300
C *EI AN ARRAY THAT ON RETURN CONTAINS THE	HQR	310
C IMAGINARY PARTS OF THE EIGENVALUES.	HQR	320
C *TYPE AN INTEGER ARRAY WHOSE I-TH ENTRY IS	HQR	330
C 0 IF THE I-TH EIGENVALUE IS REAL,	HQR	340
C 1 IF THE I-TH EIGENVALUE IS COMPLEX	HQR	350
C WITH POSITIVE IMAGINARY PART.	HQR	360
C 2 IF THE I-TH EIGENVALUE IS COMPLEX	HQR	370
C WITH NEGATIVE IMAGINARY PART,	HQR	380
C -1 IF THE I-TH EIGENVALUE WAS NOT	HQR	390
C CALCULATED SUCCESSFULLY.	HQR	400
C NA THE FIRST DIMENSION OF THE ARRAY A.	HQR	410
C NV THE FIRST DIMENSION OF THE ARRAY V.	HQR	420
C THE CONVERGENCE CRITERION EPS IS USED TO DETERMINE	HQR	430
C WHEN A SUBDIAGONAL ELEMENT OF A IS NEGLIGIBLE.	HQR	440
C SPECIFICALLY A(I+1,I) IS REGARDED AS NEGLIGIBLE	HQR	450
C IF	HQR	460
C ABS(A(I+1,I)) .LE. EPS*(ABS(A(I,I))+ABS(A(I+1,I+1))).	HQR	470
C THIS MEANS THAT THE FINAL MATRIX RETURNED BY THE	HQR	480

```

C PROGRAM WILL BE EXACTLY SIMILAR TO A + E WHERE E IS          HQR 490
C OF ORDER EPS*NORM(A), FOR ANY REASONABLY BALANCED NORM      HQR 500
C SUCH AS THE ROW-SUM NORM.                                     HQR 510
C INTERNAL VARIABLES                                           HQR 520
  INTEGER I, IT, L, MU, NL, NU                                  HQR 530
  REAL E1, E2, P, Q, R, S, T, W, X, Y, Z                      HQR 540
  LOGICAL FAIL                                                 HQR 550
C INITIALIZE.                                                 HQR 560
  DO 10 I=NLOW,NUP                                             HQR 570
    TYPE(I) = -1                                               HQR 580
  10 CONTINUE                                                  HQR 590
    T = 0.                                                      HQR 600
C MAIN LOOP. FIND AND ORDER EIGENVALUES.                       HQR 610
  NU = NUP                                                       HQR 620
  20 IF (NU.LT.NLOW) GO TO 240                                  HQR 630
    IT = 0                                                       HQR 640
C QR LOOP. FIND NEGLIGIBLE ELEMENTS AND PERFORM               HQR 650
C QR STEPS.                                                    HQR 660
  30 CONTINUE                                                  HQR 670
C SEARCH BACK FOR NEGLIGIBLE ELEMENTS.                         HQR 680
  L = NU                                                         HQR 690
  40 CONTINUE                                                  HQR 700
    IF (L.EQ.NLOW) GO TO 50                                       HQR 710
    IF (ABS(A(L,L-1)).LT.EPS*(ABS(A(L-1,L-1))+ABS(A(L,L)))) GO TO HQR 720
    * 50                                                         HQR 730
    L = L - 1                                                    HQR 740
    GO TO 40                                                       HQR 750
  50 CONTINUE                                                  HQR 760
C TEST TO SEE IF AN EIGENVALUE OR A 2X2 BLOCK                 HQR 770
C HAS BEEN FOUND.                                             HQR 780
  X = A(NU,NU)                                                  HQR 790
  IF (L.EQ.NU) GO TO 160                                         HQR 800
  Y = A(NU-1,NU-1)                                              HQR 810
  W = A(NU,NU-1)*A(NU-1,NU)                                     HQR 820
  IF (L.EQ.NU-1) GO TO 100                                       HQR 830
C TEST ITERATION COUNT. IF IT IS 30 QUIT. IF                  HQR 840
C IT IS 10 OR 20 SET UP AN AD-HOC SHIFT.                       HQR 850
  IF (IT.EQ.30) GO TO 240                                         HQR 860
  IF (IT.NE.10 .AND. IT.NE.20) GO TO 70                          HQR 870
C AD-HOC SHIFT.                                               HQR 880
  T = T + X                                                      HQR 890
  DO 60 I=NLOW,NU                                               HQR 900
    A(I,I) = A(I,I) - X                                           HQR 910
  60 CONTINUE                                                  HQR 920
  S = ABS(A(NU,NU-1)) + ABS(A(NU-1,NU-2))                      HQR 930
  X = 0.75*S                                                     HQR 940
  Y = X                                                         HQR 950
  W = -0.4375*S**2                                              HQR 960
  70 CONTINUE                                                  HQR 970
  IT = IT + 1                                                    HQR 980
C LOOK FOR TWO CONSECUTIVE SMALL SUB-DIAGONAL                 HQR 990
C ELEMENTS.                                                    HQR 1000
  NL = NU - 2                                                    HQR 1010
  80 CONTINUE                                                  HQR 1020
  Z = A(NL,NL)                                                  HQR 1030
  R = X - Z                                                      HQR 1040
  S = Y - Z                                                      HQR 1050
  P = (R*S-W)/A(NL+1,NL) + A(NL,NL+1)                          HQR 1060
  Q = A(NL+1,NL+1) - Z - R - S                                   HQR 1070
  R = A(NL+2,NL+1)                                              HQR 1080
  S = ABS(P) + ABS(Q) + ABS(R)                                   HQR 1090
  P = P/S                                                        HQR 1100
  Q = Q/S                                                        HQR 1110
  R = R/S                                                        HQR 1120
  IF (NL.EQ.L) GO TO 90                                          HQR 1130
  IF (ABS(A(NL,NL-1))*(ABS(Q)+ABS(R)).LE.EPS*ABS(P)*(ABS(A(NL-1, HQR 1140
  * NL-1))+ABS(Z)+ABS(A(NL+1,NL+1)))) GO TO 90                 HQR 1150
  NL = NL - 1                                                    HQR 1160
  GO TO 80                                                       HQR 1170
  90 CONTINUE                                                  HQR 1180
C PERFORM A QR STEP BETWEEN NL AND NU.                         HQR 1190
  CALL QRSTEP(A, V, P, Q, R, NL, NU, N, NA, NV)                 HQR 1200
  GO TO 30                                                       HQR 1210
C 2X2 BLOCK FOUND.                                            HQR 1220
  100 IF (NU.NE.NLOW+1) A(NU-1,NU-2) = 0.                      HQR 1230
  A(NU,NU) = A(NU,NU) + T                                       HQR 1240

```

```

      A(NU-1,NU-1) = A(NU-1,NU-1) + T
      TYPE(NU) = 0
      TYPE(NU-1) = 0
      MU = NU
C LOOP TO POSITION 2X2 BLOCK.
110 CONTINUE
      NL = MU - 1
C ATTEMPT TO SPLIT THE BLOCK INTO TWO REAL
C EIGENVALUES.
      CALL SPLIT(A, V, N, NL, E1, E2, NA, NV)
C IF THE SPLIT WAS SUCCESSFUL, GO AND ORDER THE
C REAL EIGENVALUES.
      IF (A(MU,MU-1).EQ.0.) GO TO 170
C TEST TO SEE IF THE BLOCK IS PROPERLY POSITIONED,
C AND IF NOT EXCHANGE IT
      IF (MU.EQ.NUP) GO TO 230
      IF (MU.EQ.NUP-1) GO TO 130
      IF (A(MU+2,MU+1).EQ.0.) GO TO 130
C THE NEXT BLOCK IS 2X2.
      IF (A(MU-1,MU-1)*A(MU,MU)-A(MU-1,MU)*A(MU,MU-1).GE.A(MU+1,
* MU+1)*A(MU+2,MU+2)-A(MU+1,MU+2)*A(MU+2,MU+1)) GO TO 230
      CALL EXCHNG(A, V, N, NL, 2, 2, EPS, FAIL, NA, NV)
      IF (.NOT.FAIL) GO TO 120
      TYPE(NL) = -1
      TYPE(NL+1) = -1
      TYPE(NL+2) = -1
      TYPE(NL+3) = -1
      GO TO 240
120 CONTINUE
      MU = MU + 2
      GO TO 150
130 CONTINUE
C THE NEXT BLOCK IS 1X1.
      IF (A(MU-1,MU-1)*A(MU,MU)-A(MU-1,MU)*A(MU,MU-1).GE.A(MU+1,
* MU+1)**2) GO TO 230
      CALL EXCHNG(A, V, N, NL, 2, 1, EPS, FAIL, NA, NV)
      IF (.NOT.FAIL) GO TO 140
      TYPE(NL) = -1
      TYPE(NL+1) = -1
      TYPE(NL+2) = -1
      GO TO 240
140 CONTINUE
      MU = MU + 1
150 CONTINUE
      GO TO 110
C SINGLE EIGENVALUE FOUND.
160 NL = 0
      A(NU,NU) = A(NU,NU) + T
      IF (NU.NE.NLOW) A(NU,NU-1) = 0.
      TYPE(NU) = 0
      MU = NU
C LOOP TO POSITION ONE OR TWO REAL EIGENVALUES.
170 CONTINUE
C POSITION THE EIGENVALUE LOCATED AT A(NL,NL).
180 CONTINUE
      IF (MU.EQ.NUP) GO TO 220
      IF (MU.EQ.NUP-1) GO TO 200
      IF (A(MU+2,MU+1).EQ.0.) GO TO 200
C THE NEXT BLOCK IS 2X2.
      IF (A(MU,MU)**2.GE.A(MU+1,MU+1)*A(MU+2,MU+2)-A(MU+1,MU+2)*
* A(MU+2,MU+1)) GO TO 230
      CALL EXCHNG(A, V, N, MU, 1, 2, EPS, FAIL, NA, NV)
      IF (.NOT.FAIL) GO TO 190
      TYPE(MU) = -1
      TYPE(MU+1) = -1
      TYPE(MU+2) = -1
      GO TO 240
190 CONTINUE
      MU = MU + 2
      GO TO 210
200 CONTINUE
C THE NEXT BLOCK IS 1X1.
      IF (ABS(A(MU,MU)).GE.ABS(A(MU+1,MU+1))) GO TO 220
      CALL EXCHNG(A, V, N, MU, 1, 1, EPS, FAIL, NA, NV)
      MU = MU + 1
210 CONTINUE

```

HQR 1250
 HQR 1260
 HQR 1270
 HQR 1280
 HQR 1290
 HQR 1300
 HQR 1310
 HQR 1320
 HQR 1330
 HQR 1340
 HQR 1350
 HQR 1360
 HQR 1370
 HQR 1380
 HQR 1390
 HQR 1400
 HQR 1410
 HQR 1420
 HQR 1430
 HQR 1440
 HQR 1450
 HQR 1460
 HQR 1470
 HQR 1480
 HQR 1490
 HQR 1500
 HQR 1510
 HQR 1520
 HQR 1530
 HQR 1540
 HQR 1550
 HQR 1560
 HQR 1570
 HQR 1580
 HQR 1590
 HQR 1600
 HQR 1610
 HQR 1620
 HQR 1630
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 HQR 1710
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 HQR 1730
 HQR 1740
 HQR 1750
 HQR 1760
 HQR 1770
 HQR 1780
 HQR 1790
 HQR 1800
 HQR 1810
 HQR 1820
 HQR 1830
 HQR 1840
 HQR 1850
 HQR 1860
 HQR 1870
 HQR 1880
 HQR 1890
 HQR 1900
 HQR 1910
 HQR 1920
 HQR 1930
 HQR 1940
 HQR 1950
 HQR 1960
 HQR 1970
 HQR 1980
 HQR 1990
 HQR 2000


```

      GO TO 180
220 CONTINUE
      MU = NL
      NL = 0
      IF (MU.NE.0) GO TO 170
C GO BACK AND GET THE NEXT EIGENVALUE.
230 CONTINUE
      NU = L - 1
      GO TO 20
C ALL THE EIGENVALUES HAVE BEEN FOUND AND ORDERED.
C COMPUTE THEIR VALUES AND TYPE.
240 IF (NU.LT.NLOW) GO TO 260
      DO 250 I=1,NU
          A(I,I) = A(I,I) + T
250 CONTINUE
260 CONTINUE
      NU = NUP
270 CONTINUE
      IF (TYPE(NU).NE.-1) GO TO 280
      NU = NU - 1
      GO TO 310
280 CONTINUE
      IF (NU.EQ.NLOW) GO TO 290
      IF (A(NU,NU-1).EQ.0.) GO TO 290
C 2X2 BLOCK.
      CALL SPLIT(A, V, N, NU-1, E1, E2, NA, NV)
      IF (A(NU,NU-1).EQ.0.) GO TO 290
      ER(NU) = E1
      EI(NU-1) = E2
      ER(NU-1) = ER(NU)
      EI(NU) = -EI(NU-1)
      TYPE(NU-1) = 1
      TYPE(NU) = 2
      NU = NU - 2
      GO TO 300
290 CONTINUE
C SINGLE ROOT.
      ER(NU) = A(NU,NU)
      EI(NU) = 0.
      NU = NU - 1
300 CONTINUE
310 CONTINUE
      IF (NU.GE.NLOW) GO TO 270
      RETURN
      END

```

HQR 2010
 HQR 2020
 HQR 2030
 HQR 2040
 HQR 2050
 HQR 2060
 HQR 2070
 HQR 2080
 HQR 2090
 HQR 2100
 HQR 2110
 HQR 2120
 HQR 2130
 HQR 2140
 HQR 2150
 HQR 2160
 HQR 2170
 HQR 2180
 HQR 2190
 HQR 2200
 HQR 2210
 HQR 2220
 HQR 2230
 HQR 2240
 HQR 2250
 HQR 2260
 HQR 2270
 HQR 2280
 HQR 2290
 HQR 2300
 HQR 2310
 HQR 2320
 HQR 2330
 HQR 2340
 HQR 2350
 HQR 2360
 HQR 2370
 HQR 2380
 HQR 2390
 HQR 2400
 HQR 2410
 HQR 2420
 HQR 2430
 HQR 2440
 HQR 2450

```

      SUBROUTINE SPLIT(A, V, N, L, E1, E2, NA, NV)
      INTEGER L, N, NA, NV
      REAL A(NA,N), V(NV,N)
C GIVEN THE UPPER HESSENBERG MATRIX A WITH A 2X2 BLOCK
C STARTING AT A(L,L), SPLIT DETERMINES IF THE
C CORRESPONDING EIGENVALUES ARE REAL OR COMPLEX. IF THEY
C ARE REAL, A ROTATION IS DETERMINED THAT REDUCES THE
C BLOCK TO UPPER TRIANGULAR FORM WITH THE EIGENVALUE
C OF LARGEST ABSOLUTE VALUE APPEARING FIRST. THE
C ROTATION IS ACCUMULATED IN V. THE EIGENVALUES (REAL
C OR COMPLEX) ARE RETURNED IN E1 AND E2. THE PARAMETERS
C IN THE CALLING SEQUENCE ARE (STARRED PARAMETERS ARE
C ALTERED BY THE SUBROUTINE)
C *A THE UPPER HESSENBERG MATRIX WHOSE 2X2
C BLOCK IS TO BE SPLIT.
C *V THE ARRAY IN WHICH THE SPLITTING TRANS-
C FORMATION IS TO BE ACCUMULATED.
C N THE ORDER OF THE MATRIX A.
C L THE POSITION OF THE 2X2 BLOCK.
C *E1 ON RETURN IF THE EIGENVALUES ARE COMPLEX
C *E2 E1 CONTAINS THEIR COMMON REAL PART AND
C E2 CONTAINS THE POSITIVE IMAGINARY PART.
C IF THE EIGENVALUES ARE REAL, E1 CONTAINS
C THE ONE LARGEST IN ABSOLUTE VALUE AND E2
C CONTAINS THE OTHER ONE.
C NA THE FIRST DIMENSION OF THE ARRAY A.
C NV THE FIRST DIMENSION OF THE ARRAY V.
C INTERNAL VARIABLES
      INTEGER I, J, L1

```

SPL 10
 SPL 20
 SPL 30
 SPL 40
 SPL 50
 SPL 60
 SPL 70
 SPL 80
 SPL 90
 SPL 100
 SPL 110
 SPL 120
 SPL 130
 SPL 140
 SPL 150
 SPL 160
 SPL 170
 SPL 180
 SPL 190
 SPL 200
 SPL 210
 SPL 220
 SPL 230
 SPL 240
 SPL 250
 SPL 260
 SPL 270
 SPL 280
 SPL 290

REAL P, Q, R, T, U, W, X, Y, Z	SPL 300
X = A(L+1,L+1)	SPL 310
Y = A(L,L)	SPL 320
W = A(L,L+1)*A(L+1,L)	SPL 330
P = (Y-X)/2.	SPL 340
Q = P**2 + W	SPL 350
IF (Q.GE.0.) GO TO 10	SPL 360
C COMPLEX EIGENVALUE.	SPL 370
E1 = P + X	SPL 380
E2 = SQRT(-Q)	SPL 390
RETURN	SPL 400
10 CONTINUE	SPL 410
C TWO REAL EIGENVALUES. SET UP TRANSFORMATION.	SPL 420
Z = SQRT(Q)	SPL 430
IF (P.LT.0.) GO TO 20	SPL 440
Z = P + Z	SPL 450
GO TO 30	SPL 460
20 CONTINUE	SPL 470
Z = P - Z	SPL 480
30 CONTINUE	SPL 490
IF (Z.EQ.0.) GO TO 40	SPL 500
R = -W/Z	SPL 510
GO TO 50	SPL 520
40 CONTINUE	SPL 530
R = 0.	SPL 540
50 CONTINUE	SPL 550
IF (ABS(X+Z).GE.ABS(X+R)) Z = R	SPL 560
Y = Y - X - Z	SPL 570
X = -Z	SPL 580
T = A(L,L+1)	SPL 590
U = A(L+1,L)	SPL 600
IF (ABS(Y)+ABS(U).LE.ABS(T)+ABS(X)) GO TO 60	SPL 610
Q = U	SPL 620
P = Y	SPL 630
GO TO 70	SPL 640
60 CONTINUE	SPL 650
Q = X	SPL 660
P = T	SPL 670
70 CONTINUE	SPL 680
R = SQRT(P**2+Q**2)	SPL 690
IF (R.GT.0.) GO TO 80	SPL 700
E1 = A(L,L)	SPL 710
E2 = A(L+1,L+1)	SPL 720
A(L+1,L) = 0.	SPL 730
RETURN	SPL 740
80 CONTINUE	SPL 750
P = P/R	SPL 760
Q = Q/R	SPL 770
C PREMULTIPLY.	SPL 780
DO 90 J=L,N	SPL 790
Z = A(L,J)	SPL 800
A(L,J) = P*Z + Q*A(L+1,J)	SPL 810
A(L+1,J) = P*A(L+1,J) - Q*Z	SPL 820
90 CONTINUE	SPL 830
C POSTMULTIPLY.	SPL 840
L1 = L + 1	SPL 850
DO 100 I=1,L1	SPL 860
Z = A(I,L)	SPL 870
A(I,L) = P*Z + Q*A(I,L+1)	SPL 880
A(I,L+1) = P*A(I,L+1) - Q*Z	SPL 890
100 CONTINUE	SPL 900
C ACCUMULATE THE TRANSFORMATION IN V.	SPL 910
DO 110 I=1,N	SPL 920
Z = V(I,L)	SPL 930
V(I,L) = P*Z + Q*V(I,L+1)	SPL 940
V(I,L+1) = P*V(I,L+1) - Q*Z	SPL 950
110 CONTINUE	SPL 960
A(L+1,L) = 0.	SPL 970
E1 = A(L,L)	SPL 980
E2 = A(L+1,L+1)	SPL 990
RETURN	SPL 1000
END	SPL 1010
SUBROUTINE (EXCHNG(A, V, N, L, B1, B2, EPS, FAIL, NA, NV)	EXC 10
INTEGER B1, B2, L, NA, NV	EXC 20

REAL A(NA,N), EPS, V(NV,N)	EXC	30
LOGICAL FAIL	EXC	40
C GIVEN THE UPPER HESSENBERG MATRIX A WITH CONSECUTIVE	EXC	50
C B1XB1 AND B2XB2 DIAGONAL BLOCKS (B1,B2 .LE. 2)	EXC	60
C STARTING AT A(L,L), EXCHNG PRODUCES A UNITARY	EXC	70
C SIMILARITY TRANSFORMATION THAT EXCHANGES THE BLOCKS	EXC	80
C ALONG WITH THEIR EIGENVALUES. THE TRANSFORMATION	EXC	90
C IS ACCUMULATED IN V. EXCHNG REQUIRES THE SUBROUTINE	EXC	100
C QRSTEP. THE PARAMETERS IN THE CALLING SEQUENCE ARE	EXC	110
C (STARRED PARAMETERS ARE ALTERED BY THE SUBROUTINE)	EXC	120
C *A THE MATRIX WHOSE BLOCKS ARE TO BE	EXC	130
C INTERCHANGED.	EXC	140
C *V THE ARRAY INTO WHICH THE TRANSFORMATIONS	EXC	150
C ARE TO BE ACCUMULATED.	EXC	160
C N THE ORDER OF THE MATRIX A.	EXC	170
C L THE POSITION OF THE BLOCKS.	EXC	180
C B1 AN INTEGER CONTAINING THE SIZE OF THE	EXC	190
C FIRST BLOCK.	EXC	200
C B2 AN INTEGER CONTAINING THE SIZE OF THE	EXC	210
C SECOND BLOCK.	EXC	220
C EPS A CONVERGENCE CRITERION (CF. HQR3).	EXC	230
C *FAIL A LOGICAL VARIABLE WHICH IS FALSE ON A	EXC	240
C NORMAL RETURN. IF THIRTY ITERATIONS WERE	EXC	250
C PERFORMED WITHOUT CONVERGENCE, FAIL IS SET	EXC	260
C TO TRUE AND THE ELEMENT	EXC	270
C A(L+B2,L+B2-1) CANNOT BE ASSUMED ZERO.	EXC	280
C NA THE FIRST DIMENSION OF THE ARRAY A.	EXC	290
C NV THE FIRST DIMENSION OF THE ARRAY V.	EXC	300
C INTERNAL VARIABLES.	EXC	310
INTEGER I, IT, J, L1, M	EXC	320
REAL P, Q, R, S, W, X, Y, Z	EXC	330
FAIL = .FALSE.	EXC	340
IF (B1.EQ.2) GO TO 70	EXC	350
IF (B2.EQ.2) GO TO 40	EXC	360
C INTERCHANGE 1X1 AND 1X1 BLOCKS.	EXC	370
L1 = L + 1	EXC	380
Q = A(L+1,L+1) - A(L,L)	EXC	390
P = A(L,L+1)	EXC	400
R = AMAX1(P,Q)	EXC	410
IF (R.EQ.0.) RETURN	EXC	420
P = P/R	EXC	430
Q = Q/R	EXC	440
R = SQRT(P**2+Q**2)	EXC	450
P = P/R	EXC	460
Q = Q/R	EXC	470
DO 10 J=L,N	EXC	480
S = P*A(L,J) + Q*A(L+1,J)	EXC	490
A(L+1,J) = P*A(L+1,J) - Q*A(L,J)	EXC	500
A(L,J) = S	EXC	510
10 CONTINUE	EXC	520
DO 20 I=1,L1	EXC	530
S = P*A(I,L) + Q*A(I,L+1)	EXC	540
A(I,L+1) = P*A(I,L+1) - Q*A(I,L)	EXC	550
A(I,L) = S	EXC	560
20 CONTINUE	EXC	570
DO 30 I=1,N	EXC	580
S = P*V(I,L) + Q*V(I,L+1)	EXC	590
V(I,L+1) = P*V(I,L+1) - Q*V(I,L)	EXC	600
V(I,L) = S	EXC	610
30 CONTINUE	EXC	620
A(L+1,L) = 0.	EXC	630
RETURN	EXC	640
40 CONTINUE	EXC	650
C INTERCHANGE 1X1 AND 2X2 BLOCKS.	EXC	660
X = A(L,L)	EXC	670
P = 1.	EXC	680
Q = 1.	EXC	690
R = 1.	EXC	700
CALL QRSTEP(A, V, P, Q, R, L, L+2, N, NA, NV)	EXC	710
IT = 0	EXC	720
50 IT = IT + 1	EXC	730
IF (IT.LE.30) GO TO 60	EXC	740
FAIL = .TRUE.	EXC	750
RETURN	EXC	760
60 CONTINUE	EXC	770

```

P = A(L,L) - X          EXC 780
Q = A(L+1,L)           EXC 790
R = 0.                 EXC 800
CALL QRSTEP(A, V, P, Q, R, L, L+2, N, NA, NV)
IF (ABS(A(L+2,L+1)).GT.EPS*(ABS(A(L+1,L+1))+ABS(A(L+2,L+2))))
* GO TO 50             EXC 810
A(L+2,L+1) = 0.       EXC 820
RETURN                EXC 830
70 CONTINUE          EXC 840
C INTERCHANGE 2X2 AND B2XB2 BLOCKS. EXC 850
M = L + 2            EXC 860
IF (B2.EQ.2) M = M + 1 EXC 870
X = A(L+1,L+1)      EXC 880
Y = A(L,L)          EXC 890
W = A(L+1,L)*A(L,L+1) EXC 900
P = 1.              EXC 910
Q = 1.              EXC 920
R = 1.              EXC 930
CALL QRSTEP(A, V, P, Q, R, L, M, N, NA, NV)
IT = 0              EXC 940
80 IT = IT + 1      EXC 950
IF (IT.LE.30) GO TO 90 EXC 960
FAIL = .TRUE.      EXC 970
RETURN              EXC 980
90 CONTINUE          EXC 990
Z = A(L,L)          EXC 1000
R = X - Z           EXC 1010
S = Y - Z           EXC 1020
P = (R*S-W)/A(L+1,L) + A(L,L+1) EXC 1030
Q = A(L+1,L+1) - Z - R - S EXC 1040
R = A(L+2,L+1)      EXC 1050
S = ABS(P) + ABS(Q) + ABS(R) EXC 1060
P = P/S             EXC 1070
Q = Q/S             EXC 1080
R = R/S             EXC 1090
CALL QRSTEP(A, V, P, Q, R, L, M, N, NA, NV)
IF (ABS(A(M-1,M-2)).GT.EPS*(ABS(A(M-1,M-1))+ABS(A(M-2,M-2))))
* GO TO 80           EXC 1100
A(M-1,M-2) = 0.    EXC 1110
RETURN              EXC 1120
END                  EXC 1130

SUBROUTINE QRSTEP(A, V, P, Q, R, NL, NU, N, NA, NV)
INTEGER N, NA, NL, NU, NV
REAL A(NA,N), P, Q, R, V(NV,N)
C QRSTEP PERFORMS ONE IMPLICIT QR STEP ON THE
C UPPER HESSENBERG MATRIX A. THE SHIFT IS DETERMINED
C BY THE NUMBERS P,Q, AND R, AND THE STEP IS APPLIED TO
C ROWS AND COLUMNS NL THROUGH NU. THE TRANSFORMATIONS
C ARE ACCUMULATED IN V. THE PARAMETERS IN THE CALLING
C SEQUENCE ARE (STARRED APARAMETERS ARE ALTERED BY THE
C SUBROUTINE)
C *A THE UPPER HESSENBERG MATRIX ON WHICH THE
C QR STEP IS TO BE PERFORMED.
C *V THE ARRAY IN WHICH THE TRANSFORMATIONS
C ARE TO BE ACCUMULATED
C *P PARAMETERS THAT DETERMINE THE SHIFT.
C *Q
C *R
C NL THE LOWER LIMIT OF THE STEP.
C NU THE UPPER LIMIT OF THE STEP.
C N THE ORDER OF THE MATRIX A.
C NA THE FIRST DIMENSION OF THE ARRAY A.
C NV THE FIRST DIMENSION OF THE ARRAY V.
C INTERNAL VARIABLES.
INTEGER I, J, K, NL2, NL3, NUM1
REAL S, X, Y, Z
LOGICAL LAST
NL2 = NL + 2
DO 10 I=NL2,NU
A(I,I-2) = 0.
10 CONTINUE
IF (NL2.EQ.NU) GO TO 30
NL3 = NL + 3

```

```

QRS 10
QRS 20
QRS 30
QRS 40
QRS 50
QRS 60
QRS 70
QRS 80
QRS 90
QRS 100
QRS 110
QRS 120
QRS 130
QRS 140
QRS 150
QRS 160
QRS 170
QRS 180
QRS 190
QRS 200
QRS 210
QRS 220
QRS 230
QRS 240
QRS 250
QRS 260
QRS 270
QRS 280
QRS 290
QRS 300
QRS 310
QRS 320

```

DO 20 I=NL3,NU	QRS 330
A(I,I-3) = 0.	QRS 340
20 CONTINUE	QRS 350
30 CONTINUE	QRS 360
NUM1 = NU - 1	QRS 370
DO 130 K=NL,NUM1	QRS 380
C DETERMINE THE TRANSFORMATION.	QRS 390
LAST = K.EQ.NUM1	QRS 400
IF (K.EQ.NL) GO TO 40	QRS 410
P = A(K,K-1)	QRS 420
Q = A(K+1,K-1)	QRS 430
R = 0.	QRS 440
IF (.NOT.LAST) R = A(K+2,K-1)	QRS 450
X = ABS(P) + ABS(Q) + ABS(R)	QRS 460
IF (X.EQ.0.) GO TO 130	QRS 470
P = P/X	QRS 480
Q = Q/X	QRS 490
R = R/X	QRS 500
40 CONTINUE	QRS 510
S = SQRT(P**2+Q**2+R**2)	QRS 520
IF (P.LT.0.) S = -S	QRS 530
IF (K.EQ.NL) GO TO 50	QRS 540
A(K,K-1) = -S*X	QRS 550
GO TO 60	QRS 560
50 CONTINUE	QRS 570
IF (NL.NE.1) A(K,K-1) = -A(K,K-1)	QRS 580
60 CONTINUE	QRS 590
P = P + S	QRS 600
X = P/S	QRS 610
Y = Q/S	QRS 620
Z = R/S	QRS 630
Q = Q/P	QRS 640
R = R/P	QRS 650
C PREMULTIPLY.	QRS 660
DO 80 J=K,N	QRS 670
P = A(K,J) + Q*A(K+1,J)	QRS 680
IF (LAST) GO TO 70	QRS 690
P = P + R*A(K+2,J)	QRS 700
A(K+2,J) = A(K+2,J) - P*Z	QRS 710
70 CONTINUE	QRS 720
A(K+1,J) = A(K+1,J) - P*Y	QRS 730
A(K,J) = A(K,J) - P*X	QRS 740
80 CONTINUE	QRS 750
C POSTMULTIPLY.	QRS 760
J = MIN0(K+3,NU)	QRS 770
DO 100 I=1,J	QRS 780
P = X*A(I,K) + Y*A(I,K+1)	QRS 790
IF (LAST) GO TO 90	QRS 800
P = P + Z*A(I,K+2)	QRS 810
A(I,K+2) = A(I,K+2) - P*R	QRS 820
90 CONTINUE	QRS 830
A(I,K+1) = A(I,K+1) - P*Q	QRS 840
A(I,K) = A(I,K) - P	QRS 850
100 CONTINUE	QRS 860
C ACCUMULATE THE TRANSFORMATION IN V.	QRS 870
DO 120 I=1,N	QRS 880
P = X*V(I,K) + Y*V(I,K+1)	QRS 890
IF (LAST) GO TO 110	QRS 900
P = P + Z*V(I,K+2)	QRS 910
V(I,K+2) = V(I,K+2) - P*R	QRS 920
110 CONTINUE	QRS 930
V(I,K+1) = V(I,K+1) - P*Q	QRS 940
V(I,K) = V(I,K) - P	QRS 950
120 CONTINUE	QRS 960
130 CONTINUE	QRS 970
RETURN	QRS 980
END	QRS 990

REMARK ON ALGORITHM 506

HQR3 and EXCHNG: Fortran Subroutines for Calculating and Ordering the Eigenvalues of a Real Upper Hessenberg Matrix [G.W. Stewart, *ACM Trans. Math. Softw.* 2, 3(Sept. 1976), 275-280]

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The following four corrections should be made to program HQR3 [1, 2]:

- (1) Line HQR 720 in [2] should have the comparison ".LE." instead of ".LT.".
- (2) Line HQR 1850 in [2] should branch to "220" instead of "230".
- (3) Line HQR 2130 should read "DO 250 I = NLOW, NU", instead of "DO 250 I = 1, NU".
- (4) Line EXC 410 should read "R = AMAX1(ABS(P), ABS(Q))" instead of "R = AMAX1(P,Q)".

The following example illustrates why the first two corrections are necessary. Apply HQR3 to the matrix

$$\begin{bmatrix} 0 & -2 & 0 & 0 & 0 & 0 \\ 3 & 5 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 4 & 0 & 0 \\ 0 & 0 & -4 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & -1 & 0 \end{bmatrix}$$

An uncorrected version of HQR3 will yield a division by zero in line HQR 1060 of [2] as a result of the test in (1). In checking for a zero subdiagonal element, the criterion is the sum of the magnitudes of the two diagonals. As the example shows, when both diagonals are zero, the .LT. test fails for an exactly zero subdiagonal element.

Without correction (2), when exchanging two upper real roots with a lower 2×2 block, only the bottom root is tested and exchanged, leaving the possibility of misordering of the top root. The example above also demonstrates this problem.

Without correction (3), the program would give the wrong diagonals in an unusual case where part of a matrix was being triangularized and the routine failed to converge.

Correction (4) is necessary when the superdiagonal element is zero between two real roots, leaving the possibility that the maximum of (P,Q) is $P = 0$ for Q negative. (If (P,Q) are both zero, no exchange is required and the program properly returns.)

REFERENCES

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2. ALGORITHM 506. *Collected Algorithms from ACM*, ACM, New York.

ALGORITHM 507

Procedures for Quintic Natural Spline Interpolation [E1]

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Key Words and Phrases: approximation, interpolation, spline, spline approximation, quintic natural spline

CR Categories: 5.13

Language: Algol W

DESCRIPTION

1. Introduction

The purpose of the procedures presented here is to determine the interpolating quintic natural spline function $S(x)$ for the set of data points (x_i, y_i) , $i = N1, N1+1, \dots, N2$, where it is assumed that $x_{N1} < x_{N1+1} < \dots < x_{N2}$. The interpolating quintic natural spline function $S(x)$ with the knots x_{N1}, \dots, x_{N2} has the following properties: (i) $S(x)$ is a polynomial of degree 5 in each interval (x_i, x_{i+1}) , $i = N1, \dots, N2-1$. (ii) $S(x)$ and its derivatives $S'(x)$, $S''(x)$, $S'''(x)$, and $S''''(x)$ are continuous in $[x_{N1}, x_{N2}]$. (iii) $S''''(x_{N1}) = S''''(x_{N2}) = S''''(x_{N1}) = S''''(x_{N2}) = 0$. (iv) $S(x_i) = y_i$, $i = N1, \dots, N2$. It is known that if $N2 > N1+1$, then there is a unique quintic natural spline function which has the properties (i)-(iv). (See, for example, Greville [3, 4].) This spline function can be represented in the form

$$S(x) = y_i + B_i t + C_i t^2 + D_i t^3 + E_i t^4 + F_i t^5 \quad (1)$$

with $t = x - x_i$ for $x_i \leq x < x_{i+1}$, $i = N1, \dots, N2-1$.

The procedure *QUINAT* computes the coefficients B_i, C_i, D_i, E_i, F_i of the quintic natural spline represented as in eq. (1) for an arbitrary set of data points (x_i, y_i) as previously specified. This procedure is much faster than the procedure *NATSPLINE* of ACM Algorithm 472 [6] with $m = 3$. An even faster procedure, *QUINEQ*, is provided for the case in which the knots x_i are known to be equidistant. In this case it is not necessary to specify the values of x_i . The representation (1) is still used, but now $t = (x - x_i)/h$, where $h = x_{i+1} - x_i$, the constant spacing of the knots.

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If at one or more of the knots x_i one also specifies the derivative y_i' , thus requiring $S'(x_i) = y_i'$, then one has to give up the condition that $S'''(x)$ be continuous at the knot x_i . If the second derivative y_i'' is also specified, thus requiring $S''(x_i) = y_i''$, then one must also give up the condition that $S'''(x)$ be continuous at x_i . *QUINAT* is designed so that it can be used in these cases with the convention that if two consecutive knots are equal, say $x_j = x_{j+1}$, then $S(x_j) = y_j$ and $S'(x_j) = y_{j+1}'$, and if three consecutive knots are equal, say $x_j = x_{j+1} = x_{j+2}$, then $S(x_j) = y_j$, $S'(x_j) = y_j'$, and $S''(x_j) = y_{j+2}''$. Thus in order to use *QUINAT* in the case that both the value y_j and the first derivative y_j' are specified at x_j , one increases the number of knots by 1 setting $x_{j+1} = x_j$ (and renumbering the knots and values to the right). Then one chooses $y_{j+1} = y_j'$. The spline function computed by *QUINAT* will have the property $S(x_j) = y_j$ and $S'(x_j) = y_{j+1}'$. One may use *QUINAT* in a similar manner if the second derivative is also specified at a knot x_j . Complete details are given in the comment of the procedure *QUINAT*.

If the values of the function y_i and the values of the first derivative y_i' are specified at all the knots x_i , then $S'''(x)$ need not be continuous at the knots and $S'''(x_{N1})$ and $S'''(x_{N2})$ need not be zero. Such a spline is said to be of deficiency 2. The procedure *QUINDF* computes the coefficients of the quintic natural spline of deficiency 2 when the values of the function y_i and the values of the first derivative y_i' are given at each knot. Although *QUINAT* could be used for this case as just described, *QUINDF* is much faster and needs much less storage space.

It is not of interest to specify the values of the function and its first and second derivatives at each knot, because in this case the quintic polynomial is completely determined in each interval independently of all other intervals.

2. Method of Calculation

QUINAT. As in the general case of Algorithm 472 [6], the calculation of the coefficients of the spline function is carried out in a numerically stable manner following a method described by Anselone and Laurent [1]. The basic ideas on which the method is based were given earlier by Schoenberg [7]. The method is specialized to the case of the quintic natural spline and uses minimum support *B*-splines [2, 3] of degree 2 to form a basis for the class of third derivatives of the quintic natural splines. Instead of specializing the formulas of Algorithm 472 [6] by setting $m = 3$, we derive the necessary formulas directly and we choose a different numbering and a different normalization for the *B*-splines.

We first assume that the knots are strictly monotone increasing, that is, $x_{N1} < x_{N1+1} < \dots < x_{N2}$. In order to simplify the notation, we choose $N1 = 0$ and let $N2 = n$, so that the data points are denoted by (x_i, y_i) , $i = 0, 1, \dots, n$. We denote by $M_i(x)$ the *B*-spline of degree 2, vanishing outside the interval (x_{i-1}, x_{i+2}) . We let $h_i = x_{i+1} - x_i$, $t = x - x_{i-1}$, $u = x - x_i$, $v = x - x_{i+1}$.

Then we have

$$\begin{aligned} M_i(x) &= At^2, & x_{i-1} \leq x < x_i, \\ &= B + Cu - Du^2, & x_i \leq x < x_{i+1}, \\ &= E(v - h_{i+1})^2, & x_{i+1} \leq x < x_{i+2}, \end{aligned} \quad (2)$$

where

$$\begin{aligned} A &= 1/[h_{i-1}(h_{i-1} + h_i)], & B &= h_{i-1}/(h_{i-1} + h_i), & C &= 2/(h_{i-1} + h_i), \\ D &= (h_{i-1} + 2h_i + h_{i+1})/[(h_{i-1} + h_i)h_i(h_i + h_{i+1})], & E &= 1/[h_{i+1}(h_i + h_{i+1})]. \end{aligned} \quad (3)$$

Now since the third derivative $S'''(x)$ vanishes outside the interval (x_0, x_n) , it has a unique representation of the form

$$S'''(x) = \sum_{j=1}^{n-2} 60\gamma_j M_j(x). \quad (4)$$

In order to determine the γ_j , we make use of the relation

$$\int_{-\infty}^{\infty} M_i(x) S'''(x) dx = 2(S(x_i, x_{i+1}, x_{i+2}) - S(x_{i-1}, x_i, x_{i+1})) \quad (5)$$

using the usual notation for divided differences. This relation is easily obtained by integration by parts. If we multiply eq. (4) by $\frac{1}{2}M_i(x)$, $i = 1, 2, \dots, n-2$, and integrate, we obtain a well-conditioned positive definite pentadiagonal system of linear equations for the determination of the γ_j :

$$\begin{aligned} d_1\gamma_1 + e_1\gamma_2 + f_1\gamma_3 &= c_1 \\ e_1\gamma_1 + d_2\gamma_2 + e_2\gamma_3 + f_2\gamma_4 &= c_2 \\ f_{i-2}\gamma_{i-2} + e_{i-1}\gamma_{i-1} + d_i\gamma_i + e_i\gamma_{i+1} + f_i\gamma_{i+2} &= c_i, \quad i = 3, 4, \dots, n-4 \\ f_{n-5}\gamma_{n-5} + e_{n-4}\gamma_{n-4} + d_{n-3}\gamma_{n-3} + e_{n-3}\gamma_{n-2} &= c_{n-3} \\ f_{n-4}\gamma_{n-4} + e_{n-3}\gamma_{n-3} + d_{n-2}\gamma_{n-2} &= c_{n-2} \end{aligned} \quad (6)$$

where

$$\begin{aligned} d_i &= T_1 + T_2 + T_3, & i &= 1, 2, \dots, n-2, \\ e_i &= T_4 + T_5, & i &= 1, 2, \dots, n-3, \\ f_i &= T_6, & i &= 1, 2, \dots, n-4, \\ c_i &= y_{i,i+1,i+2} - y_{i-1,i,i+1}, & i &= 1, 2, \dots, n-2. \end{aligned}$$

Here $y_{i,i+1,i+2}$ denotes the second divided difference of the given $\{y_i\}$, and for the T_i one finds, after some algebraic manipulation, the following formulas:

$$\begin{aligned} T_1 &= 6h_{i-1}^3/(h_{i-1} + h_i)^2 \\ T_2 &= h_i\{30h_{i-1}^2h_{i+1}^2 + (h_{i-1} + h_{i+1})h_i(40h_{i-1}h_{i+1} + 14h_i^2) \\ &\quad + h_i^2[16(h_{i-1}^2 + h_{i+1}^2) + 42h_{i-1}h_{i+1} + 4h_i^2]\}/[(h_{i-1} + h_i)^2(h_i + h_{i+1})^2] \\ T_3 &= 6h_{i+1}^3/(h_i + h_{i+1})^2 \\ T_4 &= h_i^2[h_{i-1}(h_i + h_{i+1}) + 3(h_{i-1} + h_i)(h_i + 3h_{i+1})]/[(h_{i-1} + h_i)(h_i + h_{i+1})^2] \\ T_5 &= h_{i+1}^2[h_{i+2}(h_i + h_{i+1}) + 3(h_{i+1} + h_{i+2})(3h_i + h_{i+1})]/[(h_i + h_{i+1})^2(h_{i+1} + h_{i+2})] \\ T_6 &= h_{i+1}^3/[(h_i + h_{i+1})(h_{i+1} + h_{i+2})]. \end{aligned}$$

Note that all terms in these expressions are positive; consequently no cancellations can occur. The system of equations (6) can be solved for the γ_j by using Gaussian elimination without pivoting. When the coefficients γ_j have been found, $S'''(x)$ is given by eq. (4). Remembering that $M_j(x)$ vanishes outside the interval (x_{j-1}, x_{j+2}) and making use of eqs. (1) through (4), we easily find that

$$\left. \begin{aligned} D_i/10 &= (\gamma_{i-1}h_i + \gamma_i h_{i-1})/(h_{i-1} + h_i) \\ E_i/5 &= (\gamma_i - \gamma_{i-1})/(h_{i-1} + h_i) \\ F_i &= (1/h_i)[(\gamma_{i+1} - \gamma_i)/(h_i + h_{i+1}) \\ &\quad - (\gamma_i - \gamma_{i-1})/(h_{i-1} + h_i)]. \end{aligned} \right\} \quad i = 2, 3, \dots, n-3.$$

These formulas can also be used for $i = 0, 1, n-2, n-1$ by adding the convention that $\gamma_{-1} = \gamma_0 = \gamma_{n-1} = \gamma_n = 0$. (Note that $D_0 = E_0 = 0$ as they should.) Finally we make use of the continuity of $S(x)$ and its first four derivatives at x_i to obtain the following formulas for B_i and C_i :

$$\begin{aligned} B_i &= \frac{h_{i-1}}{h_{i-1} + h_i} \frac{y_{i+1} - y_i}{h_i} + \frac{h_i}{h_{i-1} + h_i} \frac{y_i - y_{i-1}}{h_{i-1}} - D_i h_{i-1} h_i + E_i h_{i-1} h_i (h_{i-1} - h_i) \\ &\quad - \frac{h_{i-1} h_i}{h_{i-1} + h_i} (F_{i-1} h_{i-1}^3 + F_i h_i^3) \\ C_i &= \frac{1}{h_{i-1} + h_i} \left(\frac{y_{i+1} - y_i}{h_i} - \frac{y_i - y_{i-1}}{h_{i-1}} \right) + D_i (h_{i-1} - h_i) \\ &\quad - E_i \frac{h_{i-1}^3 + h_i^3}{h_{i-1} + h_i} + \frac{1}{h_{i-1} + h_i} (F_{i-1} h_{i-1}^4 - F_i h_i^4). \end{aligned}$$

These formulas are valid for $i = 1, 2, \dots, n-1$. In addition, we have for the end-points:

$$C_0 = C_1 - 10F_0 h_0^3, \quad B_0 = (y_1 - y_0)/h_0 - C_0 h_0 - F_0 h_0^4,$$

$$C_n = C_{n-1} + 10F_{n-1}h_{n-1}^3, \quad B_n = (y_n - y_{n-1})/h_{n-1} + C_n h_{n-1} - F_{n-1} h_{n-1}^4.$$

In the preceding discussion we have assumed that the knots were distinct. We can relax this condition and allow two or three consecutive knots to be equal. The procedure *QUINAT* has been written in such a way that if $x_j = x_{j+1}$, then $S(x_j) = y_j$ and $S'(x_j) = y_{j+1}$, and if $x_j = x_{j+1} = x_{j+2}$, then, in addition, $S''(x_j) = y_{j+2}$. The use of *QUINAT* in these cases is fully explained in its comment.

QUINEQ. The calculation of the coefficients in *QUINEQ* for the case of equidistant knots is carried out in the same manner as is the calculation of the coefficients in *QUINAT* for the general case. However, there are a number of simplifications which result in considerable economy of computational effort. It is not necessary to specify x_i . Hence we can assume $x_i = i$. Then $h_i = 1$ for all i , and the coefficients of $M_i(x)$ are independent of i as are also the d_i, e_i, f_i of the pentadiagonal system (6) for the γ_i . Thus eqs. (2) reduce to

$$M_i(x) = \begin{cases} \frac{1}{2} t^2, & i - 1 \leq x < i, \\ \frac{1}{2} + u - u^2, & i \leq x < i + 1, \\ \frac{1}{2} (v - 1)^2, & i + 1 \leq x < i + 2, \end{cases}$$

with $t = x - (i - 1)$, $u = x - i$, $v = x - (i + 1)$.

Instead of eq. (4) it is convenient to take

$$S'''(x) = \sum_{j=0}^{n-3} 120\gamma_j M_{j+1}(x).$$

The divided differences become ordinary differences so that eq. (5) becomes:

$$\int_{-\infty}^{\infty} M_i(x) S'''(x) dx = \Delta^3 S(x_{i-1}).$$

The pentadiagonal system (6) for the determination of γ_j becomes:

$$\begin{aligned} 66\gamma_0 + 26\gamma_1 + \gamma_2 &= \Delta^3 y_0 \\ 26\gamma_0 + 66\gamma_1 + 26\gamma_2 + \gamma_3 &= \Delta^3 y_1 \\ \gamma_{i-2} + 26\gamma_{i-1} + 66\gamma_i + 26\gamma_{i+1} + \gamma_{i+2} &= \Delta^3 y_i, \quad i = 2, 3, \dots, n-5 \\ \gamma_{n-6} + 26\gamma_{n-5} + 66\gamma_{n-4} + 26\gamma_{n-3} &= \Delta^3 y_{n-4} \\ \gamma_{n-5} + 26\gamma_{n-4} + 66\gamma_{n-3} &= \Delta^3 y_{n-3}. \end{aligned}$$

The equations for the determination of the spline function coefficients then become:

$$\begin{aligned} D_i/10 &= \gamma_{i-2} + \gamma_{i-1} & B_i &= \frac{1}{2} (y_{i+1} - y_{i-1} - F_{i-1} - F_i) - D_i \\ E_i/5 &= \gamma_{i-1} - \gamma_{i-2} & C_i &= \frac{1}{2} (y_{i+1} + y_{i-1} + F_{i-1} - F_i) - y_i - E_i. \\ F_i &= \gamma_i - \gamma_{i-1} - \gamma_{i-1} + \gamma_{i-2} \end{aligned}$$

These formulas are valid for $i = 1, 2, \dots, n-1$ with the convention that $\gamma_{-1} = \gamma_{n-2} = \gamma_{n-1} = 0$. The formula for F_i can be used for $i = 0$ by setting $\gamma_{-2} = 0$. (Note that $D_0 = E_0 = 0$ as they should.) Finally the coefficients B_i and C_i at the endpoints are given by

$$\begin{aligned} C_0 &= C_1 - 10F_0, & B_0 &= y_1 - y_0 - C_0 - F_0, \\ C_n &= C_{n+1} + 10F_{n-1}, & B_n &= y_n - y_{n-1} + C_n - F_{n-1}. \end{aligned}$$

QUINDF. We now assume $S(x_i) = y_i$ and $S'(x_i) = y_i'$ are specified at each of the knots. We must exclude the possibility that $x_i = x_{i+1}$ as this would imply a multiplicity of 4, which is not feasible for quintic splines.

We could proceed as in the calculation of *QUINAT* by using minimum support B -splines of degree 2 to form a basis for the class of third derivatives of the quintic natural splines. Of course, the B -splines would also have to be of deficiency 2. We

would again obtain a pentadiagonal system of equations which could be solved and then the coefficients for the deficient quintic natural spline could be calculated. An algorithm based on this method was developed and tested by the present authors [5].

However, we have found that we can obtain a more efficient algorithm by imposing the appropriate continuity conditions directly on eq. (1) at the knots. A similar method was used by Späth [8] to obtain an algorithm for the deficient quintic spline but with different end conditions. (He specified the second derivatives at the endpoints of the interval instead of requiring the third derivative to be zero at the endpoints as for the natural spline.)

Using eq. (1) we see at once that $S'(x_i) = y_i'$ implies $B_i = y_i'$. Then imposing the requirement that $S(x)$, $S'(x)$, and $S''(x)$ be continuous at x_{i+1} , $i = 0, 1, \dots, n-2$, we obtain by setting $t = h_i = x_{i+1} - x_i$ in eq. (1):

$$\begin{aligned} y_{i+1} &= y_i + B_i h_i + C_i h_i^2 + D_i h_i^3 + E_i h_i^4 + F_i h_i^5 \\ B_{i+1} &= B_i + 2C_i h_i + 3D_i h_i^2 + 4E_i h_i^3 + 5F_i h_i^4 \\ C_{i+1} &= C_i + 3D_i h_i + 6E_i h_i^2 + 10F_i h_i^3. \end{aligned} \quad (7)$$

If we multiply these three equations by $10/h_i^3$, $-4/h_i^2$, $1/h_i$, respectively, and add the results, we eliminate E_i and F_i and obtain

$$D_i = 10 \frac{y_{i+1} - y_i}{h_i^3} - \frac{4B_{i+1} + 6B_i}{h_i^2} + \frac{C_{i+1} - 3C_i}{h_i}. \quad (8)$$

We note that $D_i = S'''(x_i + 0)/6$. In order to obtain a similar expression for $S'''(x_i - 0)/6$, we replace the subscript $i+1$ by $i-1$ consistently (noting that then $h_i = x_{i+1} - x_i$ is to be replaced by $x_{i-1} - x_i = -h_{i-1}$). We obtain

$$\frac{1}{6} S'''(x_i - 0) = 10 \frac{y_i - y_{i-1}}{h_{i-1}^3} - \frac{6B_i + 4B_{i-1}}{h_{i-1}^2} + \frac{3C_i - C_{i-1}}{h_{i-1}}. \quad (9)$$

Now since the third derivative is continuous at x_i , we can equate the values of $S'''(x_i + 0)/6$ and $S'''(x_i - 0)/6$ in eqs. (8) and (9) to obtain the following set of equations for the C_i :

$$\begin{aligned} -\frac{1}{h_{i-1}} C_{i-1} + \left(\frac{3}{h_{i-1}} + \frac{3}{h_i} \right) C_i - \frac{1}{h_i} C_{i+1} \\ = 10 \left(\frac{y_{i+1} - y_i}{h_i^3} - \frac{y_i - y_{i-1}}{h_{i-1}^3} \right) - \frac{4B_{i+1} + 6B_i}{h_i^2} + \frac{6B_i + 4B_{i-1}}{h_{i-1}^2}. \end{aligned}$$

These equations hold for $i = 1, 2, \dots, n-1$. Two additional equations can be obtained from the conditions $S'''(x_0) = S'''(x_n) = 0$ by setting $i = 0$ in eq. (8) and by setting $i = n$ in eq. (9):

$$\begin{aligned} \frac{3}{h_0} C_0 - \frac{1}{h_0} C_1 &= 10 \frac{y_1 - y_0}{h_0^3} - \frac{4B_1 + 6B_0}{h_0^2}, \\ -\frac{1}{h_{n-1}} C_{n-1} + \frac{3}{h_n} C_n &= -10 \frac{y_n - y_{n-1}}{h_{n-1}^3} + \frac{6B_n + 4B_{n-1}}{h_{n-1}^2}. \end{aligned}$$

This system of $n+1$ equations is solved to obtain the C_i . If we make the substitutions

$$\begin{aligned} x &= D_i, & p &= (y_{i+1} - y_i - B_i h_i - C_i h_i^2)/h_i^3, \\ y &= E_i h_i, & q &= (B_{i+1} - B_i - 2C_i h_i)/h_i^2, \\ z &= F_i h_i^2, & r &= (C_{i+1} - C_i)/h_i, \end{aligned}$$

then for each i eqs. (7) form a system of three equations in the three unknowns x, y, z , and the system is solved by Gaussian elimination. The backward substitution

yields formulas for z, y, x in the following order:

$$\begin{cases} g = q - 3p \\ z = r - 3(p + g) \\ y = g - z - z \\ x = p - y - z \\ F_i = z/h_i^2 \\ E_i = y/h_i \\ D_i = x. \end{cases}$$

3. Tests

These procedures have been tested in Algol 60 on the Telefunken TR-440 computer at the Leibniz-Rechenzentrum of the Bavarian Academy of Sciences, Munich, and in Algol W on the IBM 360/67 at the Stanford Center for Information Processing. The latter tests included timing tests of the procedures with the number of knots $N = N_2 - N_1 + 1$ ranging up to 1000. The time was found to be approximately proportional to the number N of knots. The time T in seconds for the execution of the procedure *QUINAT* was found to be approximately $T = .00193N$, whereas for the procedure *NATSPLINE* of Algorithm 472 [6] with $m = 3$ it was found to be $T = .0120N$, or over six times as great. For the procedure *QUINEQ* the time was approximately $T = .00064N$, whereas for the procedure *NATSPLINEEQ* of Algorithm 472 [6] with $m = 3$ it was $T = .0038N$, or nearly six times as great. For the procedure *QUINDF* the time was approximately $T = .00087N$, whereas for the procedure *QUINAT* with $2N$ knots, consecutive knots being equal in pairs, the time was $T = .00325N$, or nearly four times as great. Moreover, to compute the same results the procedure *QUINAT* requires approximately twice as much storage for the arrays used as does the procedure *QUINDF*. Note also that from the preceding formula for the time required by the procedure *QUINAT*, the time for $2N$ distinct knots would be $T = .00386N$, which can be compared with $T = .00325N$ given above for N pairs of equal knots. The reduction for the case of double knots occurs because some calculations are omitted when knots are coincident.

These timing comparisons show that it is definitely advantageous to use these special procedures for the quintic natural spline rather than the general cases given in Algorithm 472 [6] with $m = 3$.

Tests of the accuracy and correctness of the coefficients computed by the procedures *QUINAT*, *QUINEQ*, and *QUINDF* were carried out as described in Algorithm 472 [6]. Table I shows the results of a typical run using *QUINDF* for 5 nonequidistant points. The values of the function and its first derivatives were specified. The first line of each entry gives the tabulated quantities at the given

Table I. Quintic Spline Calculated by *QUINDF*
(Machine precision approximately 7 decimal digits.)

x	S(x)	S'(x)	S''(x)/2	S'''(x)/6	S''''(x)/24	S ^v (x)/120
-3.000000	7.000000	2.000000	-6.108377	-5.722046'-06	2.956286	-0.7145951
	11.000000	15.000000	7.674876	-4.933508	-4.189662	-0.7145951
-1.000000	11.000000	15.000000	7.674870	-4.933474	-8.157658	5.416262
	26.000000	9.999996	-1.908880	16.59851	18.92365	5.416262
0	26.000000	10.000000	-1.908880	16.59848	-9.059000	1.246088
	55.99942	-27.00066	-5.264791	20.03839	9.632320	1.246088
3.000000	56.000000	-27.000000	-5.264426	20.03847	-21.28366	6.509618
	29.000000	-30.000000	-7.754811	-1.907349'-06	11.26443	6.509618
4.000000	29.000000	-30.000000				

value of x which is the lefthand endpoint of the subinterval; the second line of each entry gives the tabulated values at the righthand endpoint of the same subinterval. The close agreement of the quantities $S(x)$, $S'(x)$, $S''(x)/2$, and $S'''(x)/6$ shows that the quintic spline function and its derivatives satisfy the required continuity conditions. This is a good indication of the correctness of the results. Note that the fourth and fifth derivatives are discontinuous. Essentially the same results were obtained by using *QUINAT* with 10 knots, in equal pairs. In addition, accuracy and timing tests were carried out for large values of N , including $N = 1000$ and 5000 , and produced very satisfactory results.

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ALGORITHM

```

PROCEDURE QUINAT(INTEGER VALUE N1,N2; REAL ARRAY X,Y,B,C,D,E,F(*));      1.
COMMENT QUINAT COMPUTES THE COEFFICIENTS OF A QUINTIC NATURAL SPLINE      2.
S(X) INTERPOLATING THE ORDINATES Y(I) AT POINTS X(I), I = N1           3.
THROUGH N2. FOR XX IN (X(I),X(I+1)) THE VALUE OF THE SPLINE           4.
FUNCTION S(XX) IS GIVEN BY THE FIFTH DEGREE POLYNOMIAL:                5.
S(XX) = (((F(I)*T+E(I))*T+D(I))*T+C(I))*T+B(I))*T+Y(I)                6.
WITH T = XX - X(I).                                                    7.
INPUT:                                                                    8.
  N1,N2 SUBSCRIPT OF FIRST AND LAST DATA POINT RESPECTIVELY,          9.
  IT IS REQUIRED THAT N2 > N1 + 1,                                       10.
  X,Y(N1:N2) ARRAYS WITH X(I) AS ABSCISSA AND Y(I) AS ORDINATE         11.
  OF THE I-TH DATA POINT. THE ELEMENTS OF THE ARRAY X                 12.
  MUST BE STRICTLY MONOTONE INCREASING (BUT SEE BELOW FOR              13.
  EXCEPTIONS TO THIS).                                                  14.
OUTPUT:                                                                    15.
  B,C,D,E,F(N1:N2) ARRAYS COLLECTING THE COEFFICIENTS OF THE          16.
  QUINTIC NATURAL SPLINE S(XX) AS DESCRIBED ABOVE.                     17.
  SPECIFICALLY B(I) = S'(X(I)), C(I) = S''(X(I))/2,                   18.
  D(I) = S'''(X(I))/6, E(I) = S''''(X(I))/24,                          19.
  F(I) = S'''''(X(I)+0)/120. F(N2) IS NEITHER USED OR                 20.
  ALTERED. THE ARRAYS B,C,D,E,F MUST ALWAYS BE DISTINCT.             21.
OPTIONS:                                                                    22.
  1. THE REQUIREMENT THAT THE ELEMENTS OF THE ARRAY X BE              23.
  STRICTLY MONOTONE INCREASING CAN BE RELAXED TO ALLOW TWO             24.
  OR THREE CONSECUTIVE ABSCISSAS TO BE EQUAL AND THEN                 25.
  SPECIFYING VALUES OF THE FIRST AND SECOND DERIVATIVES OF           26.
  THE SPLINE FUNCTION AT SOME OF THE INTERPOLATING POINTS.           27.
  SPECIFICALLY                                                         28.
  IF X(J) = X(J+1) THEN S(X(J)) = Y(J) AND S'(X(J)) = Y(J+1),        29.
  IF X(J) = X(J+1) = X(J+2) THEN IN ADDITION S''(X(J)) = Y(J+2).    30.
  NOTE THAT S''''(X) IS DISCONTINUOUS AT A DOUBLE KNOT AND IN         31.
  ADDITION S''''(X) IS DISCONTINUOUS AT A TRIPLE KNOT. AT A          32.
  DOUBLE KNOT, X(J) = X(J+1), THE OUTPUT COEFFICIENTS HAVE THE        33.
  FOLLOWING VALUES:                                                    34.
  B(J) = S'(X(J)) = B(J+1)                                             35.
  C(J) = S''(X(J))/2 = C(J+1)                                          36.
  D(J) = S'''(X(J))/6 = D(J+1)                                         37.
  E(J) = S''''(X(J)-0)/24 E(J+1) = S''''(X(J)+0)/24                 38.
  F(J) = S'''''(X(J)-0)/120 F(J+1) = S'''''(X(J)+0)/120             39.
  THE REPRESENTATION OF S(XX) REMAINS VALID IN ALL INTERVALS         40.
  PROVIDED THE REDEFINITION Y(J+1) := Y(J) IS MADE                    41.

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IMMEDIATELY AFTER THE CALL OF THE PROCEDURE QUINAT. AT A 42.
TRIPLE KNOT, X(J) = X(J+1) = X(J+2), THE OUTPUT COEFFICIENTS 43.
HAVE THE FOLLOWING VALUES: 44.
  B(J) = S'(X(J)) = B(J+1) = B(J+2) 45.
  C(J) = S''(X(J))/2 = C(J+1) = C(J+2) 46.
  D(J) = S'''(X(J)-0)/6 D(J+1) = 0 D(J+2) = S'''(X(J)+0)/6 47.
  E(J) = S''''(X(J)-0)/24 E(J+1) = 0 E(J+2) = S''''(X(J)+0)/24 48.
  F(J) = S'''''(X(J)-0)/120 F(J+1)=0 F(J+2)=S'''''(X(J)+0)/120 49.
THE REPRESENTATION OF S(X) REMAINS VALID IN ALL INTERVALS 50.
PROVIDED THE REDEFINITION Y(J+2) := Y(J+1) := Y(J) IS MADE 51.
IMMEDIATELY AFTER THE CALL OF THE PROCEDURE QUINAT. 52.
2. THE ARRAY X MAY BE MONOTONE DECREASING INSTEAD OF 53.
   INCREASING; 54.
IF N2 > N1 + 1 THEN 55.
BEGIN 56.
  INTEGER M; 57.
  REAL B1,P,PQ,PQR,PR,P2,P3,Q,QR,Q2,Q3,R,R2,S,T,U,V; 58.
  COMMENT COEFFICIENTS OF A POSITIVE DEFINITE, PENTADIAGONAL 59.
  MATRIX STORED IN D,E,F(N1+1:N2-2); 60.
  M:=N2-2; 61.
  Q:=X(N1+1)-X(N1); R:=X(N1+2)-X(N1+1); 62.
  Q2:=Q*Q; R2:=R*R; QR:=Q+R; 63.
  D(N1):=E(N1):=0.0; 64.
  D(N1+1):=IF Q=0.0 THEN 0.0 ELSE 6.0*Q*Q2/(QR*QR); 65.
  FOR I:=N1+1 STEP 1 UNTIL M DO 66.
  BEGIN 67.
    P:=Q; Q:=R; R:=X(I+2)-X(I+1); 68.
    P2:=Q2; Q2:=R2; R2:=R*R; PQ:=QR; QR:=Q+R; 69.
    IF Q=0.0 THEN D(I+1):=E(I):=F(I-1):=0.0 ELSE 70.
    BEGIN 71.
      Q3:=Q2*Q; PR:=P*R; PQR:=P*Q*QR; 72.
      D(I+1):=6.0*Q3/(QR*QR); 73.
      D(I):=D(I)+(Q+Q)*(15.0*PR*PR+(P+R)*Q*(20.0*PR+7.0*Q2) 74.
        +Q2*(8.0*(P2+R2)+21.0*PR+Q2+Q2))/(PQR*PQR); 75.
      D(I-1):=D(I-1)+6.0*Q3/(PQ*PQ); 76.
      E(I):=Q2*(P*QR+3.0*PQ*(QR+R+R))/(PQR*QR); 77.
      E(I-1):=E(I-1)+Q2*(R*PQ+3.0*QR*(PQ+P+P))/(PQR*PQ); 78.
      F(I-1):=Q3/PQR; 79.
    END; 80.
  END; 81.
  IF R=0.0 THEN D(M):=D(M)+6.0*R*R2/(QR*QR); 82.
  COMMENT FIRST AND SECOND ORDER DIVIDED DIFFERENCES OF THE GIVEN 83.
  FUNCTION VALUES, STORED IN B(N1+1:N2) AND C(N1+2:N2) 84.
  RESPECTIVELY, TAKE CARE OF DOUBLE AND TRIPLE KNOTS; 85.
  S:=Y(N1); 86.
  FOR I:=N1+1 STEP 1 UNTIL N2 DO 87.
  IF X(I)=X(I-1) THEN B(I):=Y(I) ELSE 88.
  BEGIN 89.
    B(I):=(Y(I)-S)/(X(I)-X(I-1)); 90.
    S:=Y(I); 91.
  END; 92.
  FOR I:=N1+2 STEP 1 UNTIL N2 DO 93.
  IF X(I)=X(I-2) THEN 94.
  BEGIN C(I):=Y(I)*0.5; B(I):=B(I-1) END 95.
  ELSE C(I):=(B(I)-B(I-1))/(X(I)-X(I-2)); 96.
  COMMENT SOLVE THE LINEAR SYSTEM WITH C(I+2)-C(I+1) 97.
  AS RIGHT-HAND SIDE; 98.
  IF M > N1 THEN 99.
  BEGIN 100.
    P:=C(N1):=E(M):=F(N1):=F(M-1):=F(M):=0.0; 101.
    C(N1+1):=C(N1+3)-C(N1+2); D(N1+1):=1.0/D(N1+1); 102.
  END; 103.
  FOR I:=N1+2 STEP 1 UNTIL M DO 104.
  BEGIN 105.
    Q:=D(I-1)*E(I-1); 106.
    D(I):=1.0/(D(I)-P*F(I-2)-Q*E(I-1)); 107.
    E(I):=E(I)-Q*F(I-1); 108.
    C(I):=C(I+2)-C(I+1)-P*C(I-2)-Q*C(I-1); 109.
    P:=D(I-1)*F(I-1); 110.
  END; 111.
  M:=N1+1; C(N2-1):=C(N2):=0.0; 112.
  FOR I:=N2-2 STEP -1 UNTIL M DO 113.
  C(I):=(C(I)-E(I)*C(I+1)-F(I)*C(I+2))*D(I); 114.
  COMMENT INTEGRATE THE THIRD DERIVATIVE OF S(X); 115.
  M:=N2-1; 116.
  Q:=X(N1+1)-X(N1); R:=X(N1+2)-X(N1+1); B1:=B(N1+1); 117.
  Q3:=Q*Q*Q; QR:=Q+R; 118.
  V:=T:=IF QR=0.0 THEN 0.0 ELSE C(N1+1)/QR; 119.
  F(N1):=IF Q=0.0 THEN 0.0 ELSE V/Q; 120.
  FOR I:=N1+1 STEP 1 UNTIL M DO 121.
  BEGIN 122.
    P:=Q; Q:=R; 123.
    R:=IF I=N2-1 THEN 0.0 ELSE X(I+2)-X(I+1); 124.
    P3:=Q3; Q3:=Q*Q*Q; PQ:=QR; QR:=Q+R; 125.

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S:=T; T:=IF QR=0.0 THEN 0.0 ELSE (C(I+1)-C(I))/QR; 126.
U:=V; V:=T-S; 127.
IF PQ=0.0 THEN 128.
BEGIN C(I):=0.5*Y(I+1); D(I):=E(I):=F(I):=0.0 END 129.
ELSE 130.
BEGIN 131.
F(I):=IF Q=0.0 THEN F(I-1) ELSE V/Q; 132.
E(I):=5.0*S; 133.
D(I):=10.0*(C(I)-Q*S); 134.
C(I):=D(I)*(P-Q)+(B(I+1)-B(I)+(U-E(I))*P 135.
-(V+E(I))*Q3)/PQ; 136.
B(I):=(P*(B(I+1)-V*Q3)+Q*(B(I)-U*P3))/PQ 137.
-P*Q*(D(I)+E(I)*(Q-P)); 138.
END; 139.
END I; 140.
COMMENT END POINTS X(N1) AND X(N2); 141.
P:=X(N1+1)-X(N1); S:=F(N1)*P*P*P; 142.
E(N1):=D(N1):=0.0; 143.
C(N1):=C(N1+1)-10.0*S; 144.
B(N1):=B1-(C(N1)+S)*P; 145.
Q:=X(N2)-X(N2-1); T:=F(N2-1)*Q*Q*Q; 146.
E(N2):=D(N2):=0.0; 147.
C(N2):=C(N2-1)+10.0*T; 148.
B(N2):=B(N2)+(C(N2)-T)*Q; 149.
END QUINAT; 150.

PROCEDURE QUINEQ(INTEGER VALUE N1,N2; REAL ARRAY Y,B,C,D,E,F(*)); 151.
COMMENT QUINEQ COMPUTES THE COEFFICIENTS OF A QUINTIC NATURAL SPLINE 152.
S(X) INTERPOLATING THE ORDINATES Y(I) AT EQUIDISTANT POINTS X(I), 153.
I = N1 THROUGH N2. FOR XX IN (X(I),X(I+1)) THE VALUE OF THE 154.
SPLINE FUNCTION S(XX) IS GIVEN BY THE FIFTH DEGREE POLYNOMIAL: 155.
S(XX) = (((F(I)*T+E(I))*T+D(I))*T+C(I))*T+B(I))*T+Y(I) 156.
WITH T = (XX - X(I))/(X(I+1) - X(I)). 157.
INPUT: 158.
N1, N2 SUBSCRIPT OF FIRST AND LAST DATA POINT RESPECTIVELY, 159.
IT IS REQUIRED THAT N2 > N1 + 1, 160.
Y(N1):N2) THE GIVEN FUNCTION VALUES (ORDINATES). 161.
OUTPUT: 162.
B,C,D,E,F(N1:N2) ARRAYS COLLECTING THE COEFFICIENTS OF THE 163.
QUINTIC NATURAL SPLINE S(XX) AS DESCRIBED ABOVE. 164.
SPECIFICALLY B(I) = S'(X(I)), C(I) = S''(X(I))/2, 165.
D(I) = S'''(X(I))/6, E(I) = S''''(X(I))/24, 166.
F(I) = S'''''(X(I)+0)/120. F(N2) IS NEITHER USED 167.
NOR ALTERED. THE ARRAYS Y,B,C,D MUST ALWAYS BE 168.
DISTINCT. IF E AND F ARE NOT WANTED, THE CALL 169.
QUINEQ(N1,N2,Y,B,C,D,D,D) MAY BE USED TO SAVE STORAGE 170.
LOCATIONS; 171.
IF N2>N1+1 THEN 172.
BEGIN 173.
INTEGER N; 174.
REAL P,Q,R,S,T,U,V; 175.
N:=N2-3; P:=Q:=R:=S:=T:=0.0; 176.
FOR I:=N1 STEP 1 UNTIL N DO 177.
BEGIN 178.
U:=P*R; B(I):=1.0/(66.0-U*R-Q); 179.
C(I):=R:=26.0-U; 180.
D(I):=Y(I+3)-3.0*(Y(I+2)-Y(I+1))-Y(I)-U*S-Q*T; 181.
Q:=P; P:=B(I); T:=S; S:=D(I) 182.
END I; 183.
D(N+1):=D(N+2):=0.0; 184.
FOR I:=N STEP -1 UNTIL N1 DO 185.
D(I):=(D(I)-C(I)*D(I+1)-D(I+2))*B(I); 186.
N:=N2-1; Q:=0.0; R:=T:=V:=D(N1); 187.
FOR I:=N1+1 STEP 1 UNTIL N DO 188.
BEGIN 189.
P:=Q; Q:=R; R:=D(I); S:=T; 190.
F(I):=T:=P-Q-Q+R; 191.
E(I):=U:=5.0*(-P+Q); 192.
D(I):=10.0*(P+Q); 193.
C(I):=0.5*(Y(I+1)+Y(I-1)+S-T)-Y(I)-U; 194.
B(I):=0.5*(Y(I+1)-Y(I-1)-S-T)-D(I) 195.
END I; 196.
F(N1):=V; E(N1):=E(N2):=D(N1):=D(N2):=0.0; 197.
C(N1):=C(N1+1)-10.0*V; C(N2):=C(N2-1)+10.0*T; 198.
B(N1):=Y(N1+1)-Y(N1)-C(N1)-V; B(N2):=Y(N2)-Y(N2-1)+C(N2)-T 199.
END QUINEQ; 200.

PROCEDURE QUINDF(INTEGER VALUE N1,N2; REAL ARRAY X,Y,B,C,D,E,F(*)); 201.
COMMENT QUINDF COMPUTES THE COEFFICIENTS OF A QUINTIC NATURAL SPLINE 202.
S(X) FOR WHICH THE ORDINATES Y(I) AND THE FIRST DERIVATIVES B(I) 203.
ARE SPECIFIED AT POINTS X(I), I = N1 THROUGH N2. FOR XX IN 204.
(X(I),X(I+1)) THE VALUE OF THE SPLINE FUNCTION S(XX) IS GIVEN 205.

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BY THE FIFTH DEGREE POLYNOMIAL: 206.
S(XX) = (((F(I)*T+E(I))*T+D(I))*T+C(I))*T+B(I))*T+Y(I) 207.
WITH T = XX - X(I). 208.
INPUT: 209.
  N1, N2 SUBSCRIPT OF FIRST AND LAST DATA POINT RESPECTIVELY, 210.
  IT IS REQUIRED THAT N2 > N1, 211.
  X,Y,B(N1:N2) ARRAYS WITH X(I) AS ABCISSA, Y(I) AS ORDINATE 212.
  AND B(I) AS FIRST DERIVATIVE AT THE I-TH DATA POINT. 213.
  THE ELEMENTS OF THE ARRAY X MUST BE STRICTLY MONOTONE 214.
  INCREASING OR DECREASING. 215.
OUTPUT: 216.
  C,D,E,F(N1:N2) ARRAYS COLLECTING THE COEFFICIENTS OF THE 217.
  QUINTIC NATURAL SPLINE S(XX) AS DESCRIBED ABOVE. 218.
  E(N2) AND F(N2) ARE NEITHER USED NOR ALTERED. THE 219.
  ARRAYS C,D,E,F MUST ALWAYS BE DISTINCT; 220.
IF N2 > N1 THEN 221.
  BEGIN 222.
  INTEGER M2; 223.
  REAL CC,G,H,HH,H2,HH2,P,PP,Q,QQ,R,RR; 224.
  M2:=N2-1; CC:=HH:=PP:=QQ:=RR:=G:=0.0; 225.
  FOR I:=N1 STEP 1 UNTIL M2 DO 226.
  BEGIN 227.
  H:=1.0/(X(I+1)-X(I)); H2:=H*H; D(I):=3.0*(HH+H) - G*HH; 228.
  P:=(Y(I+1)-Y(I))*H*H2; Q:=(B(I+1)+B(I))*H2; 229.
  R:=(B(I+1)-B(I))*H2; 230.
  C(I):=CC:=10.0*(P-PP) - 5.0*(Q-QQ) + R + RR + G*CC; 231.
  G:=H/D(I); HH:=H; HH2:=H2; PP:=P; QQ:=Q; RR:=R 232.
  END I; 233.
  C(N2):=(-10.0*PP + 5.0*QQ + RR + G*CC)/(3.0*HH - G*HH); 234.
  FOR I:=M2 STEP -1 UNTIL N1 DO 235.
  BEGIN 236.
  D(I+1):=1.0/(X(I+1)-X(I)); C(I):=(C(I) + C(I+1)*D(I+1))/D(I) 237.
  END I; 238.
  FOR I:=N1 STEP 1 UNTIL M2 DO 239.
  BEGIN 240.
  H:=D(I+1); H2:=H*H; 241.
  P:=(Y(I+1)-Y(I))*H*H2 - B(I)*H2 - C(I)*H; 242.
  Q:=(B(I+1)-B(I))*H2 - C(I)*(H+H); 243.
  R:=(C(I+1)-C(I))*H; 244.
  G:=Q - 3.0*P; RR:=R - 3.0*(P+G); QQ:= -RR - RR + G; 245.
  F(I):=RR*H2; E(I):=QQ*H; D(I):= -RR - QQ + P 246.
  END I; 247.
  D(N2):=0.0 248.
END QUINDF; 249.

```

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REMARK ON ALGORITHM 507

Procedures for Quintic Natural Spline Interpolation

[J.G. Herriot and C.H. Reinsch, *ACM Trans. Math. Softw.* 2, 3 (June 1976), 281-289.]

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Line number 82. of Algorithm 507

IF R \neq 0.0 THEN D(M) := D(M) + 6.0*R*R2/(QR*QR)

is changed to the mathematically equivalent line

IF ABS(R) > 0.0 THEN D(M) := D(M) + 6.0*R*R2/(QR*QR)

The change is made for this reason: the character " \neq " will often not transmit correctly in the *Collected Algorithms from ACM* distribution process.

ALGORITHM 508

Matrix Bandwidth and Profile Reduction [F1]

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Key Words and Phrases: bandwidth reduction, diameter of a graph, profile reduction, sparse matrices

CR Categories: 5.14, 5.32

Language: Fortran

DESCRIPTION

Introduction

This program, REDUCE, reduces the bandwidth and profile of sparse symmetric matrices, using row and corresponding column permutations. It is a realization of the algorithm described by the authors in [4]. It was extensively tested and compared with several other programs [5] and was found to be considerably faster than the others, generally superior for bandwidth reduction and as satisfactory as any other for profile reduction.

Outline of the Method

Only an outline of the algorithm is given here; a detailed description can be found in [4]. The algorithm can best be described in terms of the *adjacency graph* G , which has the characteristic that there is an edge in G between vertices v_i and v_j if and only if $a_{ij} \neq 0$ and $i \neq j$.

Step 1. Find the endpoints of a *pseudodiameter*, i.e. a pair of vertices that are at nearly maximal distance apart. This is done by a finite, iterative process of determining a vertex that is a maximum distance away from a given vertex.

Step 2. Given pseudodiameter endpoints u and v of distance k apart, partition the set of vertices into *levels* L_1, L_2, \dots, L_k such that adjacent vertices in G are in the same or adjacent levels and such that $\max_i |L_i|$ is nearly minimized.

Step 3. Number the vertices of G , level by level, beginning at an endpoint of the pseudodiameter.

Matrix Data Structure

Sparse matrices are typically stored in some compact form which takes advantage of the sparsity. The data structure assumed here is one which is commonly used in

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bandwidth and profile schemes, e.g. [1], [2], [3], and [6]. Subroutine REDUCE accepts as input a *connection table*, C , representing the indices of the nonzero elements of the $n \times n$ matrix A . The connection table has n rows and m columns where m is the number of off-diagonal nonzero elements in the row which has a maximum number of off-diagonal nonzero elements (i.e. the maximum degree of the graph G). The entries in row i of C are the column indices of the nonzero elements in row i of the matrix A . For example, if

$$A = \begin{bmatrix} X & X & 0 & 0 & X \\ X & X & 0 & X & X \\ 0 & 0 & X & X & 0 \\ 0 & X & X & X & 0 \\ X & X & 0 & 0 & X \end{bmatrix}$$

where X represents a nonzero element, then

$$C = \begin{bmatrix} 2 & 5 & 0 \\ 1 & 4 & 5 \\ 4 & 0 & 0 \\ 2 & 3 & 0 \\ 1 & 2 & 0 \end{bmatrix}.$$

The order of indices in a row of C is immaterial. The nonzero elements of the matrix A are never needed, only their indices.

Test Results

REDUCE was tested on an IBM System/360 (model 50) computer using the Fortran IV G and H compilers and on a CDC 6600 computer using the Fortran (RUN) and Fortran extended (FTN) compilers.

In another paper [5], the authors compared the execution times, bandwidths, and profiles produced by REDUCE with those of five other programs on a wide range of problems. REDUCE typically produced the smallest bandwidths; it produced profiles which were on the average as small as those for any other program; and REDUCE was faster than all the others by at least an order of magnitude.

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ALGORITHM

SUBROUTINE REDUCE(NDSTK, NR, IOLD, RENUM, NDEG, LVL, LVLS1,	RED	10
* LVLS2, CCSTOR, IBW2, IPF2)	RED	20
C SUBROUTINE REDUCE DETERMINES A ROW AND COLUMN PERMUTATION WHICH,	RED	30
C WHEN APPLIED TO A GIVEN SPARSE MATRIX, PRODUCES A PERMUTED	RED	40
C MATRIX WITH A SMALLER BANDWIDTH AND PROFILE.	RED	50
C THE INPUT ARRAY IS A CONNECTION TABLE WHICH REPRESENTS THE	RED	60
C INDICES OF THE NONZERO ELEMENTS OF THE MATRIX, A. THE ALGO-	RED	70
C RITHM IS DESCRIBED IN TERMS OF THE ADJACENCY GRAPH WHICH	RED	80
C HAS THE CHARACTERISTIC THAT THERE IS AN EDGE (CONNECTION)	RED	90
C BETWEEN NODES I AND J IF A(I,J) .NE. 0 AND I .NE. J.	RED	100
C DIMENSIONING INFORMATION--THE FOLLOWING INTEGER ARRAYS MUST BE	RED	110
C DIMENSIONED IN THE CALLING ROUTINE.	RED	120
C NDSTK(NR,D1) D1 IS .GE. MAXIMUM DEGREE OF ALL NODES.	RED	130
C IOLD(D2) D2 AND NR ARE .GE. THE TOTAL NUMBER OF	RED	140
C RENUM(D2+1) NODES IN THE GRAPH.	RED	150

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C   NDEG(D2)          STORAGE REQUIREMENTS CAN BE SIGNIFICANTLY      RED 160
C   LVL(D2)          DECREASED FOR IBM 360 AND 370 COMPUTERS        RED 170
C   LVLS1(D2)        BY REPLACING INTEGER NDSTK BY                  RED 180
C   LVLS2(D2)        INTEGER*2 NDSTK IN SUBROUTINES REDUCE,        RED 190
C   CCSTOR(D2)       DGREE, FNDIAM, TREE AND NUMBER.              RED 200
C   COMMON INFORMATION--THE FOLLOWING COMMON BLOCK MUST BE IN THE  RED 210
C   CALLING ROUTINE.                                             RED 220
C   COMMON/GRA/N, IDPTH, IDEG                                     RED 230
C   EXPLANATION OF INPUT VARIABLES--                             RED 240
C   NDSTK-           CONNECTION TABLE REPRESENTING GRAPH.        RED 250
C                   NDSTK(I,J)=NODE NUMBER OF JTH CONNECTION TO   RED 260
C                   NUMBER I. A CONNECTION OF A NODE TO ITSELF IS NOT RED 270
C                   LISTED. EXTRA POSITIONS MUST HAVE ZERO FILL.  RED 280
C   NR-              ROW DIMENSION ASSIGNED NDSTK IN CALLING PROGRAM. RED 290
C   IOLD(I)-         NUMBERING OF ITH NODE UPON INPUT.            RED 300
C                   IF NO NUMBERING EXISTS THEN IOLD(I)=I.        RED 310
C   N-               NUMBER OF NODES IN GRAPH (EQUAL TO ORDER OF MATRIX). RED 320
C   IDEG-            MAXIMUM DEGREE OF ANY NODE IN THE GRAPH.     RED 330
C   EXPLANATION OF OUTPUT VARIABLES--                             RED 340
C   RENUM(I)-        THE NEW NUMBER FOR THE ITH NODE..           RED 350
C   NDEG(I)-         THE DEGREE OF THE ITH NODE.                 RED 360
C   IBW2-            THE BANDWIDTH AFTER RENUMBERING.            RED 370
C   IPF2-            THE PROFILE AFTER RENUMBERING.              RED 380
C   IDPTH-           NUMBER OF LEVELS IN REDUCE LEVEL STRUCTURE.  RED 390
C   THE FOLLOWING ONLY HAVE MEANING IF THE GRAPH WAS CONNECTED-- RED 400
C   LVL(I)-          INDEX INTO LVLS1 TO THE FIRST NODE IN LEVEL I. RED 410
C                   LVL(I+1)-LVL(I)= NUMBER OF NODES IN ITH LEVEL RED 420
C   LVLS1-           NODE NUMBERS LISTED BY LEVEL.              RED 430
C   LVLS2(I)-        THE LEVEL ASSIGNED TO NODE I BY REDUCE.     RED 440
C   WORKING STORAGE VARIABLE--                                   RED 450
C   CCSTOR                                                   RED 460
C   LOCAL STORAGE--                                           RED 470
C   COMMON/CC/-SUBROUTINES REDUCE, SORT2 AND PIKLV L ASSUME THAT  RED 480
C                   THE GRAPH HAS AT MOST 50 CONNECTED COMPONENTS. RED 490
C                   SUBROUTINE FNDIAM ASSUMES THAT THERE ARE AT MOST RED 500
C                   100 NODES IN THE LAST LEVEL.                 RED 510
C   COMMON/LVLW/-SUBROUTINES SETUP AND PIKLV L ASSUME THAT THERE  RED 520
C                   ARE AT MOST 100 LEVELS.                      RED 530
C   USE INTEGER*2 NDSTK WITH AN IBM 360 OR 370.                RED 540
C   INTEGER NDSTK                                             RED 550
C   INTEGER STNODE, RVNODE, RENUM, XC, SORT2, STNUM, CCSTOR,    RED 560
C   * SIZE, STPT, SBNUM                                       RED 570
C   COMMON /GRA/ N, IDPTH, IDEG                                RED 580
C   IT IS ASSUMED THAT THE GRAPH HAS AT MOST 50 CONNECTED COMPONENTS. RED 590
C   COMMON /CC/ XC, SIZE(50), STPT(50)                        RED 600
C   COMMON /LVLW/ NHIGH(100), NLOW(100), NACUM(100)           RED 610
C   DIMENSION CCSTOR(1), IOLD(1)                             RED 620
C   DIMENSION NDSTK(NR,1), LVL(1), LVLS1(1), LVLS2(1), RENUM(1), RED 630
C   * NDEG(1)                                                 RED 640
C   IBW2 = 0                                                  RED 650
C   IPF2 = 0                                                  RED 660
C   SET RENUM(I)=0 FOR ALL I TO INDICATE NODE I IS UNNUMBERED RED 670
C   DO 10 I=1,N                                              RED 680
C       RENUM(I) = 0                                         RED 690
C   10 CONTINUE                                              RED 700
C   COMPUTE DEGREE OF EACH NODE AND ORIGINAL BANDWIDTH AND PROFILE RED 710
C   CALL DGREE(NDSTK, NR, NDEG, IOLD, IBW1, IPF1)            RED 720
C   SBNUM= LOW END OF AVAILABLE NUMBERS FOR RENUMBERING      RED 730
C   STNUM= HIGH END OF AVAILABLE NUMBERS FOR RENUMBERING     RED 740
C   SBNUM = 1                                               RED 750
C   STNUM = N                                               RED 760
C   NUMBER THE NODES OF DEGREE ZERO                           RED 770
C   DO 20 I=1,N                                              RED 780
C       IF (NDEG(I).GT.0) GO TO 20                            RED 790
C       RENUM(I) = STNUM                                     RED 800
C       STNUM = STNUM - 1                                    RED 810
C   20 CONTINUE                                              RED 820
C   FIND AN UNNUMBERED NODE OF MIN DEGREE TO START ON        RED 830
C   30 LOWDG = IDEG + 1                                       RED 840
C   NFLG = 1                                               RED 850
C   ISDIR = 1                                               RED 860
C   DO 40 I=1,N                                              RED 870
C       IF (NDEG(I).GE.LOWDG) GO TO 40                       RED 880
C       IF (RENUM(I).GT.0) GO TO 40                         RED 890
C       LOWDG = NDEG(I)                                     RED 900
C       STNODE = I                                         RED 910

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40 CONTINUE                                RED 920
C FIND PSEUDO-DIAMETER AND ASSOCIATED LEVEL STRUCTURES. RED 930
C STNODE AND RVNODE ARE THE ENDS OF THE DIAM AND LVLS1 AND LVLS2 RED 940
C ARE THE RESPECTIVE LEVEL STRUCTURES. RED 950
  CALL FNDIAM(STNODE, RVNODE, NDSTK, NR, NDEG, LVL, LVLS1,
  * LVLS2, CCSTOR, IDFLT) RED 960
  IF (NDEG(STNODE).LE.NDEG(RVNODE)) GO TO 50 RED 970
C NFLG INDICATES THE END TO BEGIN NUMBERING ON RED 980
  NFLG = -1 RED 990
  STNODE = RVNODE RED 1000
50 CALL SETUP(LVL, LVLS1, LVLS2) RED 1010
C FIND ALL THE CONNECTED COMPONENTS (XC COUNTS THEM) RED 1020
  XC = 0 RED 1030
  LROOT = 1 RED 1040
  LVLN = 1 RED 1050
  DO 60 I=1,N RED 1060
    IF (LVL(I).NE.0) GO TO 60 RED 1070
    XC = XC + 1 RED 1080
    STPT(XC) = LROOT RED 1090
    CALL TREE(I, NDSTK, NR, LVL, CCSTOR, NDEG, LVLWTH, LVLBOT,
  * LVLN, MAXLW, N) RED 1100
    SIZE(XC) = LVLBOT + LVLWTH - LROOT RED 1110
    LROOT = LVLBOT + LVLWTH RED 1120
    LVLN = LROOT RED 1130
60 CONTINUE RED 1140
  IF (SORT2(DMY).EQ.0) GO TO 70 RED 1150
  CALL PIKLV(LVLS1, LVLS2, CCSTOR, IDFLT, ISDIR) RED 1160
C ON RETURN FROM PIKLV, ISDIR INDICATES THE DIRECTION THE LARGEST RED 1170
C COMPONENT FELL. ISDIR IS MODIFIED NOW TO INDICATE THE NUMBERING RED 1180
C DIRECTION. NUM IS SET TO THE PROPER VALUE FOR THIS DIRECTION. RED 1190
70 ISDIR = ISDIR*NFLG RED 1200
  NUM = SBNUM RED 1210
  IF (ISDIR.LT.0) NUM = STNUM RED 1220
  CALL NUMBER(STNODE, NUM, NDSTK, LVLS2, NDEG, RENUM, LVLS1,
  * LVL, NR, NFLG, IBW2, IPF2, CCSTOR, ISDIR) RED 1230
C UPDATE STNUM OR SBNUM AFTER NUMBERING RED 1240
  IF (ISDIR.LT.0) STNUM = NUM RED 1250
  IF (ISDIR.GT.0) SBNUM = NUM RED 1260
  IF (SBNUM.LE.STNUM) GO TO 30 RED 1270
  IF (IBW2.LE.IBW1) RETURN RED 1280
C IF ORIGINAL NUMBERING IS BETTER THAN NEW ONE, SET UP TO RETURN IT RED 1290
  DO 80 I=1,N RED 1300
    RENUM(I) = IOLD(I) RED 1310
80 CONTINUE RED 1320
  IBW2 = IBW1 RED 1330
  IPF2 = IPF1 RED 1340
  RETURN RED 1350
  END RED 1360

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SUBROUTINE DGREE(NDSTK, NR, NDEG, IOLD, IBW1, IPF1) DGR 10
C DGREE COMPUTES THE DEGREE OF EACH NODE IN NDSTK AND STORES DGR 20
C IT IN THE ARRAY NDEG. THE BANDWIDTH AND PROFILE FOR THE ORIGINAL DGR 30
C OR INPUT RENUMBERING OF THE GRAPH IS COMPUTED ALSO. DGR 40
C USE INTEGER*2 NDSTK WITH AN IBM 360 OR 370. DGR 50
  INTEGER NDSTK DGR 60
  COMMON /GRA/ N, IDPTH, IDEG DGR 70
  DIMENSION NDSTK(NR,1), NDEG(1), IOLD(1) DGR 80
  IBW1 = 0 DGR 90
  IPF1 = 0 DGR 100
  DO 40 I=1,N DGR 110
    NDEG(I) = 0 DGR 120
    IRW = 0 DGR 130
    DO 20 J=1, IDEG DGR 140
      ITST = NDSTK(I,J) DGR 150
      IF (ITST) 30, 30, 10 DGR 160
10    NDEG(I) = NDEG(I) + 1 DGR 170
      IDIF = IOLD(I) - IOLD(ITST) DGR 180
      IF (IRW.LT.IDIF) IRW = IDIF DGR 190
20    CONTINUE DGR 200
30    IPF1 = IPF1 + IRW DGR 210
      IF (IRW.GT.IBW1) IBW1 = IRW DGR 220
40 CONTINUE DGR 230
  RETURN DGR 240
  END DGR 250

```

SUBROUTINE FNDIAM(SND1, SND2, NDSTK, NR, NDEG, LVL, LVLS1,	FND	10
* LVLS2, IWK, IDFLT)	FND	20
C FNDIAM IS THE CONTROL PROCEDURE FOR FINDING THE PSEUDO-DIAMETER OF	FND	30
C NDSTK AS WELL AS THE LEVEL STRUCTURE FROM EACH END	FND	40
C SND1- ON INPUT THIS IS THE NODE NUMBER OF THE FIRST	FND	50
C ATTEMPT AT FINDING A DIAMETER. ON OUTPUT IT	FND	60
C CONTAINS THE ACTUAL NUMBER USED.	FND	70
C SND2- ON OUTPUT CONTAINS OTHER END OF DIAMETER	FND	80
C LVLS1- ARRAY CONTAINING LEVEL STRUCTURE WITH SND1 AS ROOT	FND	90
C LVLS2- ARRAY CONTAINING LEVEL STRUCTURE WITH SND2 AS ROOT	FND	100
C IDFLT- FLAG USED IN PICKING FINAL LEVEL STRUCTURE, SET	FND	110
C =1 IF WIDTH OF LVLS1 .LE. WIDTH OF LVLS2, OTHERWISE =2	FND	120
C LVL,IWK- WORKING STORAGE	FND	130
C USE INTEGER*2 NDSTK WITH AN IBM 360 OR 370.	FND	140
INTEGER NDSTK	FND	150
INTEGER FLAG, SND, SND1, SND2	FND	160
COMMON /GRA/ N, IDPTH, IDEG	FND	170
C IT IS ASSUMED THAT THE LAST LEVEL HAS AT MOST 100 NODES.	FND	180
COMMON /CC/ NDLST(100)	FND	190
DIMENSION NDSTK(NR,1), NDEG(1), LVL(1), LVLS1(1), LVLS2(1),	FND	200
* IWK(1)	FND	210
FLAG = 0	FND	220
MTW2 = N	FND	230
SND = SND1	FND	240
C ZERO LVL TO INDICATE ALL NODES ARE AVAILABLE TO TREE	FND	250
10 DO 20 I=1,N	FND	260
LVL(I) = 0	FND	270
20 CONTINUE	FND	280
LVLN = 1	FND	290
C DROP A TREE FROM SND	FND	300
CALL TREE(SND, NDSTK, NR, LVL, IWK, NDEG, LVLWTH, LVLBOT,	FND	310
* LVLN, MAXLW, MTW2)	FND	320
IF (FLAG.GE.1) GO TO 50	FND	330
FLAG = 1	FND	340
30 IDPTH = LVLN - 1	FND	350
MTW1 = MAXLW	FND	360
C COPY LEVEL STRUCTURE INTO LVLS1	FND	370
DO 40 I=1,N	FND	380
LVLS1(I) = LVL(I)	FND	390
40 CONTINUE	FND	400
NDXN = 1	FND	410
NDXL = 0	FND	420
MTW2 = N	FND	430
C SORT LAST LEVEL BY DEGREE AND STORE IN NDLST	FND	440
CALL SORTDG(NDLST, IWK(LVLBOT), NDXL, LVLWTH, NDEG)	FND	450
SND = NDLST(1)	FND	460
GO TO 10	FND	470
50 IF (IDPTH.GE.LVLN-1) GO TO 60	FND	480
C START AGAIN WITH NEW STARTING NODE	FND	490
SND1 = SND	FND	500
GO TO 30	FND	510
60 IF (MAXLW.GE.MTW2) GO TO 80	FND	520
MTW2 = MAXLW	FND	530
SND2 = SND	FND	540
C STORE NARROWEST REVERSE LEVEL STRUCTURE IN LVLS2	FND	550
DO 70 I=1,N	FND	560
LVLS2(I) = LVL(I)	FND	570
70 CONTINUE	FND	580
80 IF (NDXN.EQ.NDXL) GO TO 90	FND	590
C TRY NEXT NODE IN NDLST	FND	600
NDXN = NDXN + 1	FND	610
SND = NDLST(NDXN)	FND	620
GO TO 10	FND	630
90 IDFLT = 1	FND	640
IF (MTW2.LE.MTW1) IDFLT = 2	FND	650
RETURN	FND	660
END	FND	670
SUBROUTINE TREE(IROOT, NDSTK, NR, LVL, IWK, NDEG, LVLWTH,	TRE	10
* LVLBOT, LVLN, MAXLW, IBORT)	TRE	20
C TREE DROPS A TREE IN NDSTK FROM IROOT	TRE	30
C LVL- ARRAY INDICATING AVAILABLE NODES IN NDSTK WITH ZERO	TRE	40
C ENTRIES. TREE ENTERS LEVEL NUMBERS ASSIGNED	TRE	50

C	DURING EXECUTION OF THIS PROCEDURE	TRE	60
C	IWK-	TRE	70
C	ON OUTPUT CONTAINS NODE NUMBERS USED IN TREE	TRE	80
C	ARRANGED BY LEVELS (IWK(LVLN) CONTAINS IROOT	TRE	90
C	AND IWK(LVLBOT+LVLWTH-1) CONTAINS LAST NODE ENTERED)	TRE	100
C	LVLWTH-	TRE	110
C	ON OUTPUT CONTAINS WIDTH OF LAST LEVEL	TRE	120
C	LVLBOT-	TRE	130
C	ON OUTPUT CONTAINS INDEX INTO IWK OF FIRST	TRE	140
C	NODE IN LAST LEVEL	TRE	150
C	MAXLW-	TRE	160
C	ON OUTPUT CONTAINS THE MAXIMUM LEVEL WIDTH	TRE	170
C	LVLN-	TRE	180
C	ON INPUT THE FIRST AVAILABLE LOCATION IN IWK	TRE	190
C	USUALLY ONE BUT IF IWK IS USED TO STORE PREVIOUS	TRE	200
C	CONNECTED COMPONENTS, LVLN IS NEXT AVAILABLE LOCATION.	TRE	210
C	ON OUTPUT THE TOTAL NUMBER OF LEVELS + 1	TRE	220
C	IBORT-	TRE	230
C	INPUT PARAM WHICH TRIGGERS EARLY RETURN IF	TRE	240
C	MAXLW BECOMES .GE. IBORT	TRE	250
C	USE INTEGER*2 NDSTK WITH AN IBM 360 OR 370.	TRE	260
	INTEGER NDSTK	TRE	270
	DIMENSION NDSTK(NR,1), LVL(1), IWK(1), NDEG(1)	TRE	280
	MAXLW = 0	TRE	290
	ITOP = LVLN	TRE	300
	INOW = LVLN	TRE	310
	LVLBOT = LVLN	TRE	320
	LVLTOP = LVLN + 1	TRE	330
	LVLN = 1	TRE	340
	LVL(IROOT) = 1	TRE	350
	IWK(ITOP) = IROOT	TRE	360
10	LVLN = LVLN + 1	TRE	370
20	IWKNOV = IWK(INOW)	TRE	380
	NDROW = NDEG(IWKNOV)	TRE	390
	DO 30 J=1,NDROW	TRE	400
	ITEST = NDSTK(IWKNOV,J)	TRE	410
	IF (LVL(ITEST).NE.0) GO TO 30	TRE	420
	LVL(ITEST) = LVLN	TRE	430
	ITOP = ITOP + 1	TRE	440
	IWK(ITOP) = ITEST	TRE	450
30	CONTINUE	TRE	460
	INOW = INOW + 1	TRE	470
	IF (INOW.LT.LVLTOP) GO TO 20	TRE	480
	LVLWTH = LVLTOP - LVLBOT	TRE	490
	IF (MAXLW.LT.LVLWTH) MAXLW = LVLWTH	TRE	500
	IF (MAXLW.GE.IBORT) RETURN	TRE	510
	IF (ITOP.LT.LVLTOP) RETURN	TRE	520
	LVLBOT = INOW	TRE	530
	LVLTOP = ITOP + 1	TRE	540
	GO TO 10	TRE	550
	END	TRE	560
	SUBROUTINE SORTDG(STK1, STK2, X1, X2, NDEG)	SOR	10
C	SORTDG SORTS STK2 BY DEGREE OF THE NODE AND ADDS IT TO THE END	SOR	20
C	OF STK1 IN ORDER OF LOWEST TO HIGHEST DEGREE. X1 AND X2 ARE THE	SOR	30
C	NUMBER OF NODES IN STK1 AND STK2 RESPECTIVELY.	SOR	40
	INTEGER X1, X2, STK1, STK2, TEMP	SOR	50
	COMMON /GRA/ N, IDPTH, IDEG	SOR	60
	DIMENSION NDEG(1), STK1(1), STK2(1)	SOR	70
	IND = X2	SOR	80
10	ITEST = 0	SOR	90
	IND = IND - 1	SOR	100
	IF (IND.LT.1) GO TO 30	SOR	110
	DO 20 I=1,IND	SOR	120
	J = I + 1	SOR	130
	ISTK2 = STK2(I)	SOR	140
	JSTK2 = STK2(J)	SOR	150
	IF (NDEG(ISTK2).LE.NDEG(JSTK2)) GO TO 20	SOR	160
	ITEST = 1	SOR	170
	TEMP = STK2(I)	SOR	180
	STK2(I) = STK2(J)	SOR	190
	STK2(J) = TEMP	SOR	200
20	CONTINUE	SOR	210
	IF (ITEST.EQ.1) GO TO 10	SOR	220
30	DO 40 I=1,X2	SOR	230
	X1 = X1 + 1	SOR	240
	STK1(X1) = STK2(I)	SOR	250
40	CONTINUE	SOR	260
	RETURN	SOR	270
	END	SOR	280

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SUBROUTINE SETUP(LVL, LVLS1, LVLS2) SET 10
C SETUP COMPUTES THE REVERSE LEVELING INFO FROM LVLS2 AND STORES SET 20
C IT INTO LVLS2. NACUM(I) IS INITIALIZED TO NODES/ITH LEVEL FOR NODES SET 30
C ON THE PSEUDO-DIAMETER OF THE GRAPH. LVL IS INITIALIZED TO NON- SET 40
C ZERO FOR NODES ON THE PSEUDO-DIAM AND NODES IN A DIFFERENT SET 50
C COMPONENT OF THE GRAPH. SET 60
COMMON /GRA/ N, IDPTH, IDEG SET 70
C IT IS ASSUMED THAT THERE ARE AT MOST 100 LEVELS. SET 80
COMMON /LVLW/ NHIGH(100), NLOW(100), NACUM(100) SET 90
DIMENSION LVL(1), LVLS1(1), LVLS2(1) SET 100
DO 10 I=1, IDPTH SET 110
  NACUM(I) = 0 SET 120
10 CONTINUE SET 130
DO 30 I=1, N SET 140
  LVL(I) = 1 SET 150
  LVLS2(I) = IDPTH + 1 - LVLS2(I) SET 160
  ITEMP = LVLS2(I) SET 170
  IF (ITEMP.GT.IDPTH) GO TO 30 SET 180
  IF (ITEMP.NE.LVLS1(I)) GO TO 20 SET 190
  NACUM(ITEMP) = NACUM(ITEMP) + 1 SET 200
  GO TO 30 SET 210
20 LVL(I) = 0 SET 220
30 CONTINUE SET 230
RETURN SET 240
END SET 250

INTEGER FUNCTION SORT2(DMY) SOR 10
C SORT2 SORTS SIZE AND STPT INTO DESCENDING ORDER ACCORDING TO SOR 20
C VALUES OF SIZE. XC=NUMBER OF ENTRIES IN EACH ARRAY SOR 30
INTEGER TEMP, CCSTOR, SIZE, STPT, XC SOR 40
C IT IS ASSUMED THAT THE GRAPH HAS AT MOST 50 CONNECTED COMPONENTS. SOR 50
COMMON /CC/ XC, SIZE(50), STPT(50) SOR 60
SORT2 = 0 SOR 70
IF (XC.EQ.0) RETURN SOR 80
SORT2 = 1 SOR 90
IND = XC SOR 100
10 ITEST = 0 SOR 110
IND = IND - 1 SOR 120
IF (IND.LT.1) RETURN SOR 130
DO 20 I=1, IND SOR 140
  J = I + 1 SOR 150
  IF (SIZE(I).GE.SIZE(J)) GO TO 20 SOR 160
  ITEST = 1 SOR 170
  TEMP = SIZE(I) SOR 180
  SIZE(I) = SIZE(J) SOR 190
  SIZE(J) = TEMP SOR 200
  TEMP = STPT(I) SOR 210
  STPT(I) = STPT(J) SOR 220
  STPT(J) = TEMP SOR 230
20 CONTINUE SOR 240
IF (ITEST.EQ.1) GO TO 10 SOR 250
RETURN SOR 260
END SOR 270

SUBROUTINE PIKLV(LVLS1, LVLS2, CCSTOR, IDFLT, ISDIR) PIK 10
C PIKLV CHOOSES THE LEVEL STRUCTURE USED IN NUMBERING GRAPH PIK 20
C LVLS1- ON INPUT CONTAINS FORWARD LEVELING INFO PIK 30
C LVLS2- ON INPUT CONTAINS REVERSE LEVELING INFO PIK 40
C ON OUTPUT THE FINAL LEVEL STRUCTURE CHOSEN PIK 50
C CCSTOR- ON INPUT CONTAINS CONNECTED COMPONENT INFO PIK 60
C IDFLT- ON INPUT =1 IF WIDTH LVLS1.LE.WIDTH LVLS2, =2 OTHERWISE PIK 70
C NHIGH KEEPS TRACK OF LEVEL WIDTHS FOR HIGH NUMBERING PIK 80
C NLOW- KEEPS TRACK OF LEVEL WIDTHS FOR LOW NUMBERING PIK 90
C NACUM- KEEPS TRACK OF LEVEL WIDTHS FOR CHOSEN LEVEL STRUCTURE PIK 100
C XC- NUMBER OF CONNECTED COMPONENTS PIK 110
C SIZE(I)- SIZE OF ITH CONNECTED COMPONENT PIK 120
C STPT(I)- INDEX INTO CCSTORE OF 1ST NODE IN ITH CON COMPT PIK 130
C ISDIR- FLAG WHICH INDICATES WHICH WAY THE LARGEST CONNECTED PIK 140
C COMPONENT FELL. =+1 IF LOW AND -1 IF HIGH PIK 150
INTEGER CCSTOR, SIZE, STPT, XC, END PIK 160
COMMON /GRA/ N, IDPTH, IDEG PIK 170
C IT IS ASSUMED THAT THE GRAPH HAS AT MOST 50 COMPONENTS AND PIK 180
C THAT THERE ARE AT MOST 100 LEVELS. PIK 190

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COMMON /LVLW/ NHIGH(100), NLOW(100), NACUM(100)	PIK 200
COMMON /CC/ XC, SIZE(50), STPT(50)	PIK 210
DIMENSION LVLS1(1), LVLS2(1), CCSTOR(1)	PIK 220
C FOR EACH CONNECTED COMPONENT DO	PIK 230
DO 80 I=1, XC	PIK 240
J = STPT(I)	PIK 250
END = SIZE(I) + J - 1	PIK 260
C SET NHIGH AND NLOW EQUAL TO NACUM	PIK 270
DO 10 K=1, IDPTH	PIK 280
NHIGH(K) = NACUM(K)	PIK 290
NLOW(K) = NACUM(K)	PIK 300
10 CONTINUE	PIK 310
C UPDATE NHIGH AND NLOW FOR EACH NODE IN CONNECTED COMPONENT	PIK 320
DO 20 K=J, END	PIK 330
INODE = CCSTOR(K)	PIK 340
LVLNH = LVLS1(INODE)	PIK 350
NHIGH(LVLNH) = NHIGH(LVLNH) + 1	PIK 360
LVLNL = LVLS2(INODE)	PIK 370
NLOW(LVLNL) = NLOW(LVLNL) + 1	PIK 380
20 CONTINUE	PIK 390
MAX1 = 0	PIK 400
MAX2 = 0	PIK 410
C SET MAX1=LARGEST NEW NUMBER IN NHIGH	PIK 420
C SET MAX2=LARGEST NEW NUMBER IN NLOW	PIK 430
DO 30 K=1, IDPTH	PIK 440
IF (2*NACUM(K).EQ.NLOW(K)+NHIGH(K)) GO TO 30	PIK 450
IF (NHIGH(K).GT.MAX1) MAX1 = NHIGH(K)	PIK 460
IF (NLOW(K).GT.MAX2) MAX2 = NLOW(K)	PIK 470
30 CONTINUE	PIK 480
C SET IT= NUMBER OF LEVEL STRUCTURE TO BE USED	PIK 490
IT = 1	PIK 500
IF (MAX1.GT.MAX2) IT = 2	PIK 510
IF (MAX1.EQ.MAX2) IT = IDFLT	PIK 520
IF (IT.EQ.2) GO TO 60	PIK 530
IF (I.EQ.1) ISDIR = -1	PIK 540
C COPY LVLS1 INTO LVLS2 FOR EACH NODE IN CONNECTED COMPONENT	PIK 550
DO 40 K=J, END	PIK 560
INODE = CCSTOR(K)	PIK 570
LVLS2(INODE) = LVLS1(INODE)	PIK 580
40 CONTINUE	PIK 590
C UPDATE NACUM TO BE THE SAME AS NHIGH	PIK 600
DO 50 K=1, IDPTH	PIK 610
NACUM(K) = NHIGH(K)	PIK 620
50 CONTINUE	PIK 630
GO TO 80	PIK 640
C UPDATE NACUM TO BE THE SAME AS NLOW	PIK 650
60 DO 70 K=1, IDPTH	PIK 660
NACUM(K) = NLOW(K)	PIK 670
70 CONTINUE	PIK 680
80 CONTINUE	PIK 690
RETURN	PIK 700
END	PIK 710
SUBROUTINE NUMBER(SND, NUM, NDSTK, LVLS2, NDEG, RENUM, LVLST,	NUM 10
* LSTPT, NR, NFLG, IBW2, IPF2, IPFA, ISDIR)	NUM 20
C NUMBER PRODUCES THE NUMBERING OF THE GRAPH FOR MIN BANDWIDTH	NUM 30
C SND- ON INPUT THE NODE TO BEGIN NUMBERING ON	NUM 40
C NUM- ON INPUT AND OUTPUT, THE NEXT AVAILABLE NUMBER	NUM 50
C LVLS2- THE LEVEL STRUCTURE TO BE USED IN NUMBERING	NUM 60
C RENUM- THE ARRAY USED TO STORE THE NEW NUMBERING	NUM 70
C LVLST- ON OUTPUT CONTAINS LEVEL STRUCTURE	NUM 80
C LSTPT(I)- ON OUTPUT, INDEX INTO LVLST TO FIRST NODE IN ITH LVL	NUM 90
C LSTPT(I+1) - LSTPT(I) = NUMBER OF NODES IN ITH LVL	NUM 100
C NFLG- =+1 IF SND IS FORWARD END OF PSEUDO-DIAM	NUM 110
C =-1 IF SND IS REVERSE END OF PSEUDO-DIAM	NUM 120
C IBW2- BANDWIDTH OF NEW NUMBERING COMPUTED BY NUMBER	NUM 130
C IPF2- PROFILE OF NEW NUMBERING COMPUTED BY NUMBER	NUM 140
C IPFA- WORKING STORAGE USED TO COMPUTE PROFILE AND BANDWIDTH	NUM 150
C ISDIR- INDICATES STEP DIRECTION USED IN NUMBERING(+1 OR -1)	NUM 160
C USE INTEGER*2 NDSTK WITH AN IBM 360 OR 370.	NUM 170
INTEGER NDSTK	NUM 180
INTEGER SND, STKA, STKB, STKC, STKD, XA, XB, XC, XD, CX, END,	NUM 190
* RENUM, TEST	NUM 200
COMMON /GRA/ N, IDPTH, IDEG	NUM 210
C THE STORAGE IN COMMON BLOCKS CC AND LVLW IS NOW FREE AND CAN	NUM 220

C BE USED FOR STACKS.	NUM 230
COMMON /LVLW/ STKA(100), STKB(100), STKC(100)	NUM 240
COMMON /CC/ STKD(100)	NUM 250
DIMENSION IPFA(1)	NUM 260
DIMENSION NDSTK(NR,1), LVLS2(1), NDEG(1), RENUM(1), LVLST(1),	NUM 270
* LSTPT(1)	NUM 280
C SET UP LVLST AND LSTPT FROM LVLS2	NUM 290
DO 10 I=1,N	NUM 300
IPFA(I) = 0	NUM 310
10 CONTINUE	NUM 320
NSTPT = 1	NUM 330
DO 30 I=1,IDPTH	NUM 340
LSTPT(I) = NSTPT	NUM 350
DO 20 J=1,N	NUM 360
IF (LVLS2(J).NE.I) GO TO 20	NUM 370
LVLST(NSTPT) = J	NUM 380
NSTPT = NSTPT + 1	NUM 390
20 CONTINUE	NUM 400
30 CONTINUE	NUM 410
LSTPT(IDPTH+1) = NSTPT	NUM 420
C STKA, STKB, STKC AND STKD ARE STACKS WITH POINTERS	NUM 430
C XA,XB,XC, AND XD. CX IS A SPECIAL POINTER INTO STKC WHICH	NUM 440
C INDICATES THE PARTICULAR NODE BEING PROCESSED.	NUM 450
C LVLN KEEPS TRACK OF THE LEVEL WE ARE WORKING AT.	NUM 460
C INITIALLY STKC CONTAINS ONLY THE INITIAL NODE, SND.	NUM 470
LVLN = 0	NUM 480
IF (NFLG.LT.0) LVLN = IDPTH + 1	NUM 490
XC = 1	NUM 500
STKC(XC) = SND	NUM 510
40 CX = 1	NUM 520
XD = 0	NUM 530
LVLN = LVLN + NFLG	NUM 540
LST = LSTPT(LVLN)	NUM 550
LND = LSTPT(LVLN+1) - 1	NUM 560
C BEGIN PROCESSING NODE STKC(CX)	NUM 570
50 IPRO = STKC(CX)	NUM 580
RENUM(IPRO) = NUM	NUM 590
NUM = NUM + ISDIR	NUM 600
END = NDEG(IPRO)	NUM 610
XA = 0	NUM 620
XB = 0	NUM 630
C CHECK ALL ADJACENT NODES	NUM 640
DO 80 I=1,END	NUM 650
TEST = NDSTK(IPRO,I)	NUM 660
INX = RENUM(TEST)	NUM 670
C ONLY NODES NOT NUMBERED OR ALREADY ON A STACK ARE ADDED	NUM 680
IF (INX.EQ.0) GO TO 60	NUM 690
IF (INX.LT.0) GO TO 80	NUM 700
C DO PRELIMINARY BANDWIDTH AND PROFILE CALCULATIONS	NUM 710
NBW = (RENUM(IPRO)-INX)*ISDIR	NUM 720
IF (ISDIR.GT.0) INX = RENUM(IPRO)	NUM 730
IF (IPFA(INX).LT.NBW) IPFA(INX) = NBW	NUM 740
GO TO 80	NUM 750
60 RENUM(TEST) = -1	NUM 760
C PUT NODES ON SAME LEVEL ON STKA, ALL OTHERS ON STKB	NUM 770
IF (LVLS2(TEST).EQ.LVLS2(IPRO)) GO TO 70	NUM 780
XB = XB + 1	NUM 790
STKB(XB) = TEST	NUM 800
GO TO 80	NUM 810
70 XA = XA + 1	NUM 820
STKA(XA) = TEST	NUM 830
80 CONTINUE	NUM 840
C SORT STKA AND STKB INTO INCREASING DEGREE AND ADD STKA TO STKC	NUM 850
C AND STKB TO STKD	NUM 860
IF (XA.EQ.0) GO TO 100	NUM 870
IF (XA.EQ.1) GO TO 90	NUM 880
CALL SORTDG(STKC, STKA, XC, XA, NDEG)	NUM 890
GO TO 100	NUM 900
90 XC = XC + 1	NUM 910
STKC(XC) = STKA(XA)	NUM 920
100 IF (XB.EQ.0) GO TO 120	NUM 930
IF (XB.EQ.1) GO TO 110	NUM 940
CALL SORTDG(STKD, STKB, XD, XB, NDEG)	NUM 950
GO TO 120	NUM 960
110 XD = XD + 1	NUM 970
STKD(XD) = STKB(XB)	NUM 980

C BE SURE TO PROCESS ALL NODES IN STKC	NUM 990
120 CX = CX + 1	NUM 1000
IF (XC.GE.CX) GO TO 50	NUM 1010
C WHEN STKC IS EXHAUSTED LOOK FOR MIN DEGREE NODE IN SAME LEVEL	NUM 1020
C WHICH HAS NOT BEEN PROCESSED	NUM 1030
MAX = IDEG + 1	NUM 1040
SND = N + 1	NUM 1050
DO 130 I=LST,LND	NUM 1060
TEST = LVLST(I)	NUM 1070
IF (RENUM(TEST).NE.0) GO TO 130	NUM 1080
IF (NDEG(TEST).GE.MAX) GO TO 130	NUM 1090
RENUM(SND) = 0	NUM 1100
RENUM(TEST) = -1	NUM 1110
MAX = NDEG(TEST)	NUM 1120
SND = TEST	NUM 1130
130 CONTINUE	NUM 1140
IF (SND.EQ.N+1) GO TO 140	NUM 1150
XC = XC + 1	NUM 1160
STKC(XC) = SND	NUM 1170
GO TO 50	NUM 1180
C IF STKD IS EMPTY WE ARE DONE, OTHERWISE COPY STKD ONTO STKC	NUM 1190
C AND BEGIN PROCESSING NEW STKC	NUM 1200
140 IF (XD.EQ.0) GO TO 160	NUM 1210
DO 150 I=1,XD	NUM 1220
STKC(I) = STKD(I)	NUM 1230
150 CONTINUE	NUM 1240
XC = XD	NUM 1250
GO TO 40	NUM 1260
C DO FINAL BANDWIDTH AND PROFILE CALCULATIONS	NUM 1270
160 DO 170 I=1,N	NUM 1280
IF (IPFA(I).GT.IBW2) IBW2 = IPFA(I)	NUM 1290
IPF2 = IPF2 + IPFA(I)	NUM 1300
170 CONTINUE	NUM 1310
RETURN	NUM 1320
END	NUM 1330

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REMARK ON ALGORITHMS 508 AND 509

Matrix Bandwidth and Profile Reduction [H.L. Crane, Jr., N.E. Gibbs, W.G. Poole, Jr., and P.K. Stockmeyer, *ACM Trans. Math. Softw.* 2, 4 (Dec. 1976), 375-377] and A Hybrid Profile Reduction Algorithm [N.E. Gibbs, *ACM Trans. Math. Softw.* 2, 4 (Dec. 1976), 378-387]

John G. Lewis [Received 28 May 1980; revised 6 January 1982; accepted 10 January 1982]

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A new implementation of the Gibbs-Poole-Stockmeyer and Gibbs-King algorithms is available as Algorithm 582. This new implementation is faster, more robust, and requires less storage than the previous implementation of the Gibbs-Poole-Stockmeyer algorithm [1], distributed as Algorithm 508, and of the Gibbs-King algorithm [2], distributed as Algorithm 509. The mathematical capabilities of Algorithm 582 are identical to those of Algorithms 508 and 509. References [3] and [4] give the implementation details and documentation for the new implementation.

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ALGORITHM 509

A Hybrid Profile Reduction Algorithm [FI]

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Key Words and Phrases: bandwidth reduction, graph, King algorithm, leveling, profile reduction, sparse matrix

CR Categories: 5.14, 5.32

Language: Fortran

DESCRIPTION

1. Introduction

Large sparse matrix problems arise in many applications areas, such as structural engineering, fluid dynamics, and network analysis. Many algorithms for automatically reordering the rows and columns of the matrix in order to reduce bandwidth and/or profile have been published. Recently, Gibbs et al. [4] concluded that of the many reduction algorithms available, theirs (as described in [3] and implemented as a Fortran program in [1]) and King's (as described in [5] and implemented as a Fortran program as described in [4]) were superior for profile reduction. Although the King algorithm produced the best profile in many cases, it sometimes exhibited erratic behavior due to the strategy of reducing the profile globally by reducing the profile well locally. Furthermore, they found a large class of problems (wide cylinders—Type E3 in [4]) for which there is no good single starting point for the King algorithm.

In this paper a new "hybrid" algorithm for reducing profile is presented. This algorithm combines the better features of [3] and [5] by first finding a pseudodiameter to produce a leveling and then numbering the leveling level by level according to King's criteria.

Section 2 describes the new profile reduction algorithm. In Section 3 it is compared in production mode with the better algorithms reported on in [4]. A listing of the Fortran subprogram used to implement the new algorithm is given in Appendix A. The changes which must be made to the Fortran subroutine REDUCE [1] to implement the algorithm are given in Appendix B.

2. Description of the Algorithm

Let $Ax = b$ be an $n \times n$ sparse nonsingular system of linear algebraic equations. To define the profile of A , first define $f_i = \min \{j : a_{ij} \neq 0\}$ for $i = 1, 2, \dots, n$ (it is assumed that $a_{ii} \neq 0$). This locates the leftmost nonzero element in each row.

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Now set $\delta_i = i - f_i$. The *profile* is defined to be $\sum_{i=1}^n \delta_i$. Given such a matrix A , we define a *graph* $G = \langle V, E \rangle$ where V has n vertices, $\{v_1, v_2, \dots, v_n\}$ and $\{v_i, v_j\} \in E$ if $a_{ij} \neq 0$ and $i \neq j$. It is well known (see, for example [3]) that the problem of matrix profile reduction can be restated as a problem of graph profile reduction. The algorithm presented here is described in terms of graphs and is directly applicable to graphs.

We assume that the reader is familiar with Algorithms I and II of Gibbs et al. [3] and with King's algorithm [5]. The reader is referred to [3] for the definitions and notation used below.

Algorithms I and II accept as input a connected graph and produce as output a leveling. A *leveling* of a graph $G = \langle V, E \rangle$ is a partition of the vertices V into levels L_1, L_2, \dots, L_k such that

- (1) all vertices adjacent to vertices in level L_1 are either in level L_1 or level L_2 ;
- (2) all vertices adjacent to vertices in level L_k are either in level L_{k-1} or level L_k ;
- (3) for $1 < i < k$, all vertices adjacent to vertices in level L_i are in either level L_{i-1} , level L_i , or level L_{i+1} .

The new algorithm, Algorithm P, described below is applied to the leveling of a connected component of a graph. The following assumptions are made.

- (1) L_1, L_2, \dots, L_k is a leveling produced by algorithms I and II and the number of levels, k , must be at least two.
- (2) On input the variable *nextnum* is initialized to the next available integer to be assigned to this connected component. On output *nextnum* will be set to the next available integer to be assigned to the next connected component.
- (3) The variables S_2, S_3 , and Q represent queues of integers. If X is a queue of integers and x is an integer variable, then:
 - (a) $X \leftarrow X - \{x\}$ means remove the integer x from the queue X (treating X as a set).
 - (b) $X \leftarrow x$ means place the integer x at the rear of the queue X .
 - (c) $x \leftarrow X$ means set the variable x to the integer at the front of the queue X after removing it from X .
 - (d) $X \leftarrow \emptyset$ means empty the queue X .
 - (e) $|X|$ represents the current size of the queue X .
- (4) Algorithm P produces as output a vector *new*, where $new(i) = j$ means that the old vertex labeled with the integer i should be relabeled with the integer j .

ALGORITHM P. PROFILE REDUCTION

- P0. [Initialization.] Set $S_2 \leftarrow L_1$ and $i \leftarrow 1$.
- P1. [Form the next level and empty Q .] Set $S_3 \leftarrow L_{i+1}$ and $Q \leftarrow \emptyset$.
- P2. [Find the next vertex to be numbered according to King's criteria.] Set m equal to the first element of S_2 which has fewest connections to S_3 . Set C equal to the set of elements of S_3 which are joined by an edge to m .
- P3. [Numbering.] Set $new(m) \leftarrow nextnum$, $S_2 \leftarrow S_2 - \{m\}$, $nextnum \leftarrow nextnum + 1$, and for each element $c \in C$ set $Q \leftarrow c$ and $S_3 \leftarrow S_3 - \{c\}$.
- P4. [Is S_2 empty?] If $|S_2| = 0$ goto step P7.
- P5. [S_2 is not empty. Is S_3 empty?] If $|S_3| > 0$ return to step P2.
- P6. [S_2 is not empty, but S_3 is empty.] For $j = 1, 2, \dots, |S_2|$ set $s \leftarrow S_2$, $new(s) \leftarrow nextnum$, and $nextnum \leftarrow nextnum + 1$. Goto step P8.
- P7. [S_2 is empty. If S_3 is not empty place the rest of S_3 in Q .] If $|S_3| > 0$ then set $q \leftarrow S_3$, $Q \leftarrow q$, and repeat step P7.
- P8. [S_2 and S_3 are empty.] Set $i \leftarrow i + 1$. If $i < k$ (the last level) then set $S_2 \leftarrow Q$ and return to step P1.
- P9. [All levels are numbered except the last level which is in Q .] If $|Q| = 0$ then terminate the algorithm; otherwise set $q \leftarrow Q$, $new(q) \leftarrow nextnum$, $nextnum \leftarrow nextnum + 1$, and repeat step P9.

Algorithm P assigns consecutive integers, level by level, to the leveling according to King's criteria. In fact, this algorithm is the King algorithm restricted to assigning consecutive integers to the vertices in one level at a time. If initially S_2 was a single vertex v , S_3 was set to $V - \{v\}$, and the elements of C were placed directly into S_2 instead of Q in step P3 we would, in essence, have King's algorithm.

It turns out that reversing the numbering new by setting $new(j) \leftarrow n - new(j) + 1$ for $j = 1, 2, \dots, n$ sometimes improves profile. We have not been able to predict a priori when reversing helps, so every numbering is reversed and the numbering that gives smaller profile is chosen (see subroutine CHECK in Appendix A.)

3. Test Results in Production Mode

The new algorithm as described above and implemented as a Fortran subroutine named PROFIT (see Appendix A) was integrated into the Fortran package described in [1]. A parameter was added to subroutine REDUCE so that the user can specify whether profile reduction is more important than bandwidth reduction or if bandwidth reduction is more important than profile reduction. The changes which were made to REDUCE are described in Appendix B.

As stated in [3], every practical bandwidth or profile reduction algorithm is heuristic in the sense that one cannot make absolute a priori statements concerning the performance of the algorithm; the performance is data dependent. In order to relate this work to previous work the new algorithm was tested on the same test data described in [4]. Tables I and II summarize the results. These tables show how subroutine REDUCE with subroutine PROFIT performed relative to Gibbs et al. (GPS [1, 3]) and King (KNG [5, 4]), with respect to bandwidth reduction, profile reduction, and time. The times are relative to 1.00 for GPS.

4. Conclusions

Although the new algorithm takes more time than GPS, it is clear that it is superior to GPS in reducing profile. The new algorithm yields profile comparable to KNG, but takes less time and does not exhibit KNG's erratic behavior.

In retrospect, the good performance of GPS for bandwidth reduction and the new algorithm for profile reduction can be explained. The Cuthill-McKee [2] and King algorithms do a good job of reducing bandwidth and profile respectively by numbering well locally. Since these algorithms examine the graph's structure locally, they sometimes behave erratically. The algorithms for producing a leveling by first finding a pseudodiameter and then using it to minimize level width take the entire structure of the graph into account and thus, when the Cuthill-McKee or King numbering is imposed on a leveling one obtains comparable bandwidth or profile respectively. The GPS algorithm and the new algorithm give superior bandwidth or profile reduction on those cases where global strategy is necessary.

Table I. Applications Problems
 See [3] for a description.
 n: order of matrix; β : bandwidth; ρ : profile; RATIO: ratio to GPS time

	n	Original		GPS		Profit			KNG		
		β	ρ	β	ρ	β	ρ	Ratio	β	ρ	Ratio
1	68	45	598	7	269	8	270	1.41	30	216	21.48
2	90	85	1020	7	579	7	579	1.52	11	573	5.66
3	92	80	2127	13	736	17	725	2.69	38	673	13.16
4	130	126	3615	18	1588	21	1482	3.43	25	1496	5.17
5	159	19	1046	12	971	15	964	1.75	19	1046	3.68
6	174	16	1569	13	1466	16	1405	2.28	28	1341	5.30
7	185	168	7534	29	3610	47	3485	6.81	116	2882	11.78
8	220	166	8532	12	1868	12	1791	2.01	13	1754	92.62
9	263	262	2681	19	2346	21	2265	2.20	28	2152	33.32
10	263	30	2040	14	2001	18	1932	1.97	30	1823	13.71
11	310	302	23357	14	2726	21	2697	2.10	25	2660	18.90
12	312	262	18076	37	5548	51	5238	7.71	190	6679	8.36
13	346	216	16435	46	7650	55	7334	10.79	66	6302	20.20
14	360	344	29790	34	6364	52	5800	6.61	316	8875	9.34
15	436	173	7913	33	7844	47	7550	4.99	173	7913	6.20
16	512	399	35837	29	4669	29	4304	1.78	59	4134	2.96
17	555	480	56322	91	28976	117	27247	59.50	179	23106	28.03
18	861	833	100560	71	45525	115	41904	37.66	436	54159	44.34
19	918	840	124607	49	20369	64	19573	5.87	737	54104	17.24

Table II. Selected Grid Problems

See [4] for a complete description of grids. * indicates test not run.

n: order of matrix; β : bandwidth; ρ : profile; RATIO: ratio to GPS timeA1: square grid (5-point difference scheme); B3: 3×1 rectangle (9-point difference scheme)D3: 20×1 cylinder (9-point difference scheme); E2: 1×20 cylinder (7-point difference scheme)

	n	Original		GPS		Profit			KNG		
		β	ρ	β	ρ	β	ρ	Ratio	β	ρ	Ratio
A1	49	47	406	7	245	7	245	1.750	7	245	27.33
A1	100	98	1153	10	705	10	705	2.111	11	705	51.57
A1	400	398	8713	20	5510	20	5510	4.068	21	5510	156.28
A1	625	623	16768	25	10700	25	10700	5.253	25	10700	158.45
B3	48	46	366	5	212	5	212	1.267	7	211	4.68
B3	432	430	7078	13	5436	13	5436	2.389	23	5371	7.02
B3	1200	1198	29406	21	24740	21	24740	3.987	39	24503	7.86
B3	1728	1726	49306	25	42552	25	42552	4.801	47	42209	8.27
D3	500	498	71474	7	3175	7	3175	1.379	9	3172	9.64
D3	980	978	218512	9	8351	9	8351	1.575	13	8341	13.19
D3	1620	1618	490830	11	17199	11	17199	1.767	17	17178	18.65
D3	2420	2418	926828	13	30679	13	30679	2.005	21	30643	22.96
E2	320	318	3367	10	2662	9	2512	2.071	13	2507	93.76
E2	720	718	10541	14	8942	14	8486	3.040	23	8468	117.30
E2	1280	1278	23931	18	21038	19	20120	4.197	*	*	*
E2	2000	1998	45457	22	40838	24	39302	5.174	*	*	*

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APPENDIX A

```

SUBROUTINE PROFIT(NR, NDSTK, NEW, NDEG, LVLS2, LVLST, LSTPT, PRO 10
* NXTNUM) PRO 20
C SUBROUTINE PROFIT NUMBERS LEVEL BY LEVEL WITH CONSECUTIVE INTEGERS PRO 30
C USING A MODIFIED VERSION OF KING*S ALGORITHM. PRO 40
C NR- ROW DIMENSION OF CONNECTION TABLE. PRO 50
C NDSTK- THE CONNECTION TABLE. PRO 60
C NEW- VECTOR TO STORE THE NEW NUMBERING. PRO 70
C NDEG(I)- THE DEGREE OF NODE I. PRO 80
C LVLS2- THE LEVEL STRUCTURE PRODUCED BY PIKLV. PRO 90
C LVLS2(I) = J MEANS VERTEX I HAS BEEN PRO 100
C PLACED IN LEVEL J. PRO 110
C LVLST- ON OUTPUT, CONTAINS THE LEVEL STRUCTURE USED. PRO 120
C LVLST(LSTPT(I)), ..., LVLST(LSTPT(I+1)-1) ARE PRO 130
C THE VERTICES IN LEVEL I. PRO 140
C LSTPT(I)- ON OUTPUT, INDEX INTO LVLST TO FIRST NODE IN LEVEL I. PRO 150
C LSTPT(I+1) - LSTPT(I) = NUMBER OF NODES IN I*TH LEVEL. PRO 160
C NXTNUM- ON INPUT AND OUTPUT, THE NEXT AVAILABLE NUMBER. PRO 170
C ON IBM 360 OR 370 USE INTEGER * 2 NDSTK. PRO 180
C INTEGER NDSTK PRO 190
C INTEGER NEW(1), NDEG(1), LVLS2(1), LVLST(1), LSTPT(1) PRO 200
C DIMENSION NDSTK(NR,1) PRO 210
C COMMON AREA GRA HOLDS VITAL INFORMATION ABOUT THE GRAPH PRO 220
C N- THE NUMBER OF NODES PRO 230
C IDPTH- THE NUMBER OF LEVELS FOUND BY PIKLV. PRO 240
C IDEG- MAXIMUM DEGREE OF GRAPH -- COLUMN DIMENSION OF NDSTK. PRO 250
C COMMON /GRA/ N, IDPTH, IDEG PRO 260
C IT IS ASSUMED THAT NO LEVEL HAS MORE THAN 100 NODES. PRO 270
C COMMON /LVLW/ S2(100), S3(100), Q(100) PRO 280
C COMMON /CC/ CONECT(100) PRO 290
C INTEGER S2, S3, Q, CONECT, S2SZE, S3SZE, QPTR, CONSZE PRO 300
C SET UP LVLST AND LSTPT FROM LVLS2. PRO 310
C NSTPT = 1 PRO 320
C DO 20 I=1, IDPTH PRO 330
C LSTPT(I) = NSTPT PRO 340
C DO 10 J=1, N PRO 350
C IF (LVLS2(J).NE.I) GO TO 10 PRO 360
C LVLST(NSTPT) = J PRO 370
C NSTPT = NSTPT + 1 PRO 380
C 10 CONTINUE PRO 390
C 20 CONTINUE PRO 400
C LSTPT(IDPTH+1) = NSTPT PRO 410
C ***** STEP P0 ***** PRO 420
C S2 IS THE FIRST LEVEL. PRO 430
C LEVEL = 1 PRO 440
C CALL FORMLV(S2, S2SZE, LSTPT, LVLST, LEVEL) PRO 450
C ***** STEP P1 ***** PRO 460
C S3 IS THE LEVEL ADJACENT TO THE LEVEL S2. PRO 470
C Q IS A QUEUE USED TO RETAIN THE ORDER IN WHICH THE ELEMENTS OF S3 PRO 480
C ARE REMOVED. Q EVENTUALLY BECOMES THE NEW S2 AND IS ORDERED PRO 490
C ACCORDING TO KING*S CRITERA. PRO 500
C 30 CALL FORMLV(S3, S3SZE, LSTPT, LVLST, LEVEL+1) PRO 510
C QPTR = 0 PRO 520
C ***** STEP P2 ***** PRO 530
C FIND THE NODE M IN S2 WHICH IS ADJACENT TO THE FEWEST NODES IN S3. PRO 540
C 40 M = MINCON(S2, S2SZE, S3, S3SZE, CONECT, CONSZE, NDSTK, NR, NDEG) PRO 550
C ***** STEP P3 ***** PRO 560
C NUMBER M AND REMOVE IT FROM S2. PRO 570
C NEW(M) = NXTNUM PRO 580
C NXTNUM = NXTNUM + 1 PRO 590

```



```

        CALL DELETE(S2, S2SZE, M)                                PRO 600
        IF (CONSZE.LE.0) GO TO 60                               PRO 610
C THE ELEMENTS OF CONLST ARE TO BE REMOVED FROM S3 AND PLACED INTO PRO 620
C Q.                                                           PRO 630
        DO 50 I=1,CONSZE                                       PRO 640
            QPTR = QPTR + 1                                     PRO 650
            Q(QPTR) = CONECT(I)                                PRO 660
            CALL DELETE(S3, S3SZE, CONECT(I))                  PRO 670
        50 CONTINUE                                           PRO 680
C ***** STEP P4 ***** PRO 690
        60 IF (S2SZE.LE.0) GO TO 80                             PRO 700
C ***** STEP P5 ***** PRO 710
        IF (S3SZE.GT.0) GO TO 40                               PRO 720
C ***** STEP P6 ***** PRO 730
C S3 IS EMPTY, BUT S2 IS NOT.  RENUMBER THE NODES WHICH REMAIN IN S2. PRO 740
        DO 70 I=1,S2SZE                                       PRO 750
            NS2 = S2(I)                                       PRO 760
            NEW(NS2) = NXTNUM                                  PRO 770
            NXTNUM = NXTNUM + 1                               PRO 780
        70 CONTINUE                                           PRO 790
        GO TO 100                                              PRO 800
C ***** STEP P7 ***** PRO 810
        80 IF (S3SZE.LE.0) GO TO 100                           PRO 820
C S2 IS EMPTY, BUT S3 IS NOT.  MOVE S3*S REMAINING NODES INTO Q. PRO 830
        DO 90 I=1,S3SZE                                       PRO 840
            QPTR = QPTR + 1                                     PRO 850
            Q(QPTR) = S3(I)                                    PRO 860
        90 CONTINUE                                           PRO 870
C ***** STEP P8 ***** PRO 880
        100 LEVEL = LEVEL + 1                                   PRO 890
            IF (LEVEL.GE.IDPTH) GO TO 120                       PRO 900
C S2 BECOMES THE OLD Q SINCE BOTH S2 AND S3 ARE EMPTY.        PRO 910
        DO 110 I=1,QPTR                                         PRO 920
            S2(I) = Q(I)                                       PRO 930
        110 CONTINUE                                           PRO 940
            S2SZE = QPTR                                       PRO 950
            GO TO 30                                           PRO 960
C ***** STEP P9 ***** PRO 970
C LAST LEVEL IS ORDERED IN Q,  SO NUMBER IT BEFORE RETURNING. PRO 980
        120 DO 130 I=1,QPTR                                     PRO 990
            IQ = Q(I)                                          PRO 1000
            NEW(IQ) = NXTNUM                                    PRO 1010
            NXTNUM = NXTNUM + 1                                PRO 1020
        130 CONTINUE                                           PRO 1030
            RETURN                                             PRO 1040
            END                                               PRO 1050

        FUNCTION MINCON(X, XSZE, Y, YSZE, CONLST, CONSZE, NDSTK, NR, MIN 10
        * NDEG)                                               MIN 20
C FUNCTION MINCON RETURNS AS ITS FUNCTIONAL VALUE A VERTEX X(I) SUCH MIN 30
C THAT THE NUMBER OF CONNECTIONS FROM X(I) TO THE SET Y IS A MINIMUM. MIN 40
C THE VERTICES OF Y WHICH ARE ADJACENT TO X(I) ARE PLACED IN MIN 50
C CONLST(1), CONLST(2),...,CONLST(CONSZE).                  MIN 60
C USE INTEGER * 2 NDSTK ON IBM 360 OR 370.                   MIN 70
        INTEGER NDSTK                                         MIN 80
        DIMENSION NDSTK(NR,1)                                  MIN 90
        INTEGER X(1), XSZE, Y(1), YSZE, CONLST(1), CONSZE, NDEG(1) MIN 100
C IT IS ASSUMED THAT NO LEVEL HAS MORE THAN 100 VERTICES.    MIN 110
        INTEGER SMLST(100)                                     MIN 120
        CONSZE = YSZE + 1                                     MIN 130
        DO 50 I=1,XSZE                                        MIN 140
            LSTSZE = 0                                         MIN 150
            IX = X(I)                                          MIN 160
            IROWDG = NDEG(IX)                                  MIN 170
            DO 20 J=1,YSZE                                     MIN 180
                DO 10 K=1,IROWDG                               MIN 190
                    IX = X(I)                                  MIN 200
                    IF (NDSTK(IX,K).NE.Y(J)) GO TO 10         MIN 210
                    SMLST(LSTSZE+1) = Y(J)                   MIN 220
                    LSTSZE = LSTSZE + 1                       MIN 230
                    IF (LSTSZE.GE.CONSZE) GO TO 50            MIN 240
                GO TO 20                                       MIN 250
            10 CONTINUE                                        MIN 260
        20 CONTINUE                                           MIN 270
            IF (LSTSZE.GT.0) GO TO 30                           MIN 280

```

C WE HAVE FOUND A VERTEX IN X WHICH IS NOT CONNECTED TO ANY VERTEX	MIN	290
C IN Y	MIN	300
MINCON = X(I)	MIN	310
CONSZE = 0	MIN	320
RETURN	MIN	330
C WE HAVE FOUND A VERTEX X(I) WITH FEWEST CONNECTIONS (NONZERO) TO Y	MIN	340
C SO FAR. SAVE THE ELEMENTS OF Y WHICH CONNECT TO X(I) IN CONLST AND	MIN	350
C SAVE X(I) AS THE FUNCTIONAL VALUE.	MIN	360
30 CONSZE = LSTSZE	MIN	370
DO 40 J=1,LSTSZE	MIN	380
CONLST(J) = SMLST(J)	MIN	390
40 CONTINUE	MIN	400
MINCON = X(I)	MIN	410
50 CONTINUE	MIN	420
RETURN	MIN	430
END	MIN	440
SUBROUTINE DELETE(SET, SETSZE, ELEMNT)	DEL	10
C SUBROUTINE DELETE REMOVES ELEMNT FROM THE SET SET IF ELEMNT	DEL	20
C IS IN SET. OTHERWISE, IT ISSUES A DIAGNOSTIC.	DEL	30
INTEGER SET(1), SETSZE, ELEMNT	DEL	40
IF (SETSZE.GT.1) GO TO 10	DEL	50
IF (SETSZE.EQ.1 .AND. SET(1).NE.ELEMNT) GO TO 30	DEL	60
SETSZE = 0	DEL	70
RETURN	DEL	80
10 DO 20 I=1,SETSZE	DEL	90
IF (SET(I).EQ.ELEMNT) GO TO 40	DEL	100
20 CONTINUE	DEL	110
30 WRITE (6,99999) ELEMNT, (SET(I),I=1,SETSZE)	DEL	120
RETURN	DEL	130
40 SETSZE = SETSZE - 1	DEL	140
DO 50 J=I,SETSZE	DEL	150
SET(J) = SET(J+1)	DEL	160
50 CONTINUE	DEL	170
RETURN	DEL	180
99999 FORMAT (10H0ERROR -- , I6, 8H NOT IN , (20I5))	DEL	190
END	DEL	200
SUBROUTINE FORMLV(SET, SETSZE, LSTPT, LVLST, LEVEL)	FOR	10
C FORMLV COPIES LEVEL(LEVEL) INTO SET.	FOR	20
INTEGER SET(1), SETSZE, LSTPT(1), LVLST(1), UPPER	FOR	30
LOWER = LSTPT(LEVEL)	FOR	40
UPPER = LSTPT(LEVEL+1) - 1	FOR	50
SETSZE = 1	FOR	60
DO 10 I=LOWER,UPPER	FOR	70
SET(SETSZE) = LVLST(I)	FOR	80
SETSZE = SETSZE + 1	FOR	90
10 CONTINUE	FOR	100
SETSZE = SETSZE - 1	FOR	110
RETURN	FOR	120
END	FOR	130
SUBROUTINE CHECK(BESTBW, BESTPF, RENUM, NDSTK, NR, NDEG, IWK)	CHE	10
C SUBROUTINE CHECK TESTS TO SEE IF REVERSED NUMBERING GIVES BETTER	CHE	20
C PROFILE THAN PROFIT. IF IT DOES, THEN RENUM IS REVERSED AND BESTPF	CHE	30
C IS SET TO THE SMALLEST OF RENUM AND REVERSED RENUM.	CHE	40
C USE INTEGER * 2 NDSTK ON IBM 360 OR 370	CHE	50
INTEGER NDSTK	CHE	60
DIMENSION NDSTK(NR,1)	CHE	70
INTEGER BESTBW, BESTPF, RENUM(1), NDEG(1), IWK(1)	CHE	80
COMMON /GRA/ N, IDPTH, IDEG	CHE	90
DO 10 I=1,N	CHE	100
IWK(I) = N - RENUM(I) + 1	CHE	110
10 CONTINUE	CHE	120
CALL BAND(BESTBW, BESTPF, RENUM, NDSTK, NR, NDEG)	CHE	130
CALL BAND(IBW, IPF, IWK, NDSTK, NR, NDEG)	CHE	140
IF (IPF.GE.BESTPF) RETURN	CHE	150
DO 20 I=1,N	CHE	160
RENUM(I) = IWK(I)	CHE	170
20 CONTINUE	CHE	180
BESTPF = IPF	CHE	190
RETURN	CHE	200
END	CHE	210

SUBROUTINE BAND(IBW, IPF, NEW, NDSTK, NR, NDEG)	BAN	10
C SUBROUTINE BAND COMPUTES THE BANDWIDTH IBW AND THE PROFILE	BAN	20
C IPF OF THE GRAPH REPRESENTED BY NDSTK USING THE NUMBERING NEW.	BAN	30
C ON IBM 360 OR 370 USE INTEGER * 2 NDSTK.	BAN	40
INTEGER NDSTK	BAN	50
DIMENSION NDSTK(NR,1), NEW(1), NDEG(1)	BAN	60
COMMON /GRA/ N, IDPTH, IDEG	BAN	70
IPF = 0	BAN	80
IBW = 0	BAN	90
DO 20 K=1,N	BAN	100
IEND = NDEG(K)	BAN	110
IF (IEND.EQ.0) GO TO 20	BAN	120
NBW = 0	BAN	130
DO 10 J=1,IEND	BAN	140
IDUMMY = NDSTK(K,J)	BAN	150
NTEST = NEW(K) - NEW(IDUMMY)	BAN	160
IF (NTEST.LE.NBW) GO TO 10	BAN	170
NBW = NTEST	BAN	180
10 CONTINUE	BAN	190
IPF = IPF + NBW	BAN	200
IF (NBW.GT.IBW) IBW = NBW	BAN	210
20 CONTINUE	BAN	220
RETURN	BAN	230
END	BAN	240

APPENDIX B. SUMMARY OF CHANGES TO FORTRAN SUBROUTINE REDUCE NECESSARY TO IMPLEMENT ALGORITHM P

There are seven changes which have to be made to REDUCE to use the subprograms given in Appendix A. The first is to add the variable OPTPRO to REDUCE's argument list. One change, the sixth, requires adding a statement number. The remaining five changes require inserting statements in the proper place in REDUCE. The update cards are underlined and at least one line of REDUCE before and after each alteration is given.

SUBROUTINE REDUCE(NDSTK, NR, IOLD, RENUM, NDEG, LVL, LVLS1,	RED	10
* LVLS2, CCSTOR, IBW2, IPF2, OPTPRO)	RED	20
C SUBROUTINE REDUCE DETERMINES A ROW AND COLUMN PERMUTATION WHICH,	RED	30
.		
C IDEG- MAXIMUM DEGREE OF ANY NODE IN THE GRAPH.	RED	330
C OPTPRO- SET TO 1 BY CALLING PROGRAM IF PROFILE REDUCTION	RED	331
<u>C IS MORE IMPORTANT THAN BANDWIDTH REDUCTION.</u>	RED	332
<u>C SET TO ANY OTHER INTEGER IF BANDWIDTH REDUCTION IS</u>	RED	333
<u>C MORE IMPORTANT THAN PROFILE REDUCTION.</u>	RED	334
C EXPLANATION OF OUTPUT VARIABLES--	RED	340
.		
INTEGER STNODE, RVNODE, RENUM, XC, SORT2, STNUM, CCSTOR,	RED	560
* SIZE, STPT, SBNUM	RED	570
INTEGER OPTPRO	RED	571
COMMON /GRA/ N, IDPTH, IDEG	RED	580
.		
CALL FNDIAM(STNODE, RVNODE, NDSTK, NR, NDEG, LVL, LVLS1,	RED	960
* LVLS2, CCSTOR, IDFLT)	RED	970
IF (OPTPRO.EQ.1) GO TO 50	RED	971
IF (NDEG(STNODE).LE.NDEG(RVNODE)) GO TO 50	RED	980
.		
70 ISDIR = ISDIR*NFLG	RED	1220
NUM = SBNUM	RED	1230
C IF PROFILE REDUCTION IS MORE IMPORTANT CALL PROFIT INSTEAD OF	RED	1231
<u>C NUMBER AND IGNORE ISDIR.</u>	RED	1232
IF (OPTPRO.EQ.1) GO TO 90	RED	1233
IF (ISDIR.LT.0) NUM = STNUM	RED	1240
.		
C IF ORIGINAL NUMBERING IS BETTER THAN NEW ONE, SET UP TO RETURN IT	RED	1320
75 DO 80 I=1,N	RED	1330
RENUM(I) = IOLD(I)	RED	1340
.		
RETURN	RED	1380

90 CALL PROFIT(NR, NDSTK, RENUM, NDEG, LVLS2, LVLS1, LVL, NUM)	RED 1381
SBNUM = NUM	RED 1382
IF (SBNUM.LE.STNUM) GO TO 30	RED 1383
C SOMETIMES PROFILE REDUCTION IS IMPROVED BY REVERSING THE NUMBERING	RED 1384
C PRODUCED BY PROFIT. SUBROUTINE CHECK DETERMINES IF THIS IS THE CASE.	RED 1385
CALL CHECK(IBW2, IPF2, RENUM, NDSTK, NR, NDEG, LVL)	RED 1386
C IF ORIGINAL NUMBERING GIVES BETTER PROFILE USE IT RATHER THAN RENUM.	RED 1387
IF (IPF2.LE.IPF1) RETURN	RED 1388
GO TO 75	RED 1389
END*	RED 1390

ACM Transactions on Mathematical Software, Vol. 8, No. 2, June 1982, Page 221.

REMARK ON ALGORITHMS 508 AND 509

Matrix Bandwidth and Profile Reduction [H.L. Crane, Jr., N.E. Gibbs, W.G. Poole, Jr., and P.K. Stockmeyer, *ACM Trans. Math. Softw.* 2, 4 (Dec. 1976), 375-377] and A Hybrid Profile Reduction Algorithm [N.E. Gibbs, *ACM Trans. Math. Softw.* 2, 4 (Dec. 1976), 378-387]

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A new implementation of the Gibbs-Poole-Stockmeyer and Gibbs-King algorithms is available as Algorithm 582. This new implementation is faster, more robust, and requires less storage than the previous implementation of the Gibbs-Poole-Stockmeyer algorithm [1], distributed as Algorithm 508, and of the Gibbs-King algorithm [2], distributed as Algorithm 509. The mathematical capabilities of Algorithm 582 are identical to those of Algorithms 508 and 509. References [3] and [4] give the implementation details and documentation for the new implementation.

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1. CRANE, H.L., JR., GIBBS, N.E., POOLE, W. G., JR., AND STOCKMEYER, P.K. Matrix bandwidth and profile reduction. *ACM Trans. Math. Softw.* 2, 4 (Dec. 1976), 375-377.
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3. LEWIS, J.G. Implementation of the Gibbs-Poole-Stockmeyer and Gibbs-King algorithms. *ACM Trans. Math. Softw.* 8, 2 (June 1982), 180-189.
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Algorithm 510

Piecewise Linear Approximations to Tabulated Data [E2]

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Key Words and Phrases: approximation, first degree spline, linear, piecewise linear
CR Categories: 5.13
Language: Fortran

DESCRIPTION

This algorithm generates a piecewise linear approximation to tabulated data which is within specified tolerances of the data points. There are eight options available. The approximation may be required to be continuous or not, a single tolerance or a tolerance for each data point may be specified, and the approximation may or may not be restricted to the piecewise linear "tolerance band" centered at the data points with edges determined by the tolerances at the data points. Among approximations of the kind requested, the algorithm gives one of fewest line segments.

The algorithm processes data in order of increasing abscissa values, successively finding the longest approximating segment by a systematic search technique. This algorithm is a descendent of the one given in [1] and the systematic search technique is the same. The specified options determine where a new segment begins and whether or not successive segments intersect.

Even if a continuous approximation is not requested, the algorithm determines whether or not successive segments would intersect at an acceptable point if the partition point between them were adjusted. If this is the case, the adjustment is made. This means that the algorithm may give a continuous approximation when one was not requested even in cases where the algorithm of [1] would not. However it does not mean that the algorithm will necessarily give a continuous approximation when one is not requested even though there may exist a continuous approximation of the same number of line segments.

In many cases there is considerable freedom of choice in determining the last segment or two of the approximation. The criteria used to make these choices are continuity first and nearness of approach to the data second.

The algorithm will accept any nonnegative tolerances. If all tolerances are zero, then the approximation linearly interpolates the data.

When the approximation is required to be continuous and not restricted to the tolerance band, the algorithm gives a first degree spline which approximates the data within the specified tolerances and is of fewest line segments.

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The inputs and outputs including the form of the approximation generated are described after the calling sequence, which is:

```
CALL STL2(X,Y,E,M,U,V,W,K,IP)
```

X, Y, E, and M contain input data; U, V, K, and possibly W contain output; and IP contains a parameter which specifies the options desired.

X and Y are real arrays of M elements. X(I) and Y(I) contain the abscissa and ordinate respectively of the I-th data point. M is an integer variable containing the dimension of the arrays X and Y.

E may be a single real variable or an M-element array depending on the value of IP. E contains the nonnegative tolerance or tolerances. If E is an array, then E(I) is the tolerance associated with (X(I), Y(I)) for $I = 1, \dots, M$.

U and V are real arrays of at least $K + 1$ elements. (Since K, the number of approximating segments, is not known a priori, the dimension of U and V must be "sufficiently large.") After STL2 has executed, U will contain a partition of the interval $[X(1), X(N)]$ with $U(1) = X(1)$ and $U(K + 1) = X(N)$. V(I) is an ordinate to be associated with U(I) in the approximation. If a continuous approximation is requested, then V(I) is "the" ordinate to be associated with U(I).

If a continuous approximation is requested, then W is not used. In this case the I-th approximating segment is the straight line from (U(I), V(I)) to (U(I + 1), V(I + 1)).

If a continuous approximation is not requested, then W is a real array of at least K elements. In this case the I-th approximating segment is the straight line from (U(I), W(I)) to (U(I + 1), V(I + 1)), and V(1) is set equal to W(1).

K is an integer variable. After STL2 has executed, K will contain either the number of approximating line segments or zero. K will be set to zero in case of an error return.

Call STL2(X,Y,E,M,X,Y,E,M,IP) will not cause problems provided that either a continuous approximation is requested, or E is a sufficiently large real array.

IP is an integer variable. The content of IP is the product of three indicators: I1, I2, and I3. I1 indicates whether or not E is an array of tolerances. $I1 = -1$ indicates that E is an array. $I1 = +1$ indicates that E is a single variable. I2 indicates whether or not the approximation is to be restricted to the tolerance band about the data. $I2 = 1$ indicates that the approximation is *not* to be restricted to the tolerance band. $I2 = 2$ indicates that the restriction *is* to be so restricted. I3 indicates whether or not the approximation is required to be continuous. $I3 = 1$ indicates that the approximation *need not be* continuous. $I3 = 3$ indicates that the approximation *must be* continuous.

The program performs the following data checks: Are the abscissa values stored in order of increasing magnitude? Is the tolerance (or are the tolerances) nonnegative? Is the number of data points greater than 1? If the answer to any of these questions is no, then the program returns with K set to zero. In this case no further processing is attempted.

TEST RESULTS

This algorithm has successfully approximated 10 sets of data specifically designed to test various parts of the program. The main routine was written to repeat each data set with the ordinate values replaced by their negatives. This provides a new approximation problem for the algorithm, essentially doubling the number of data sets, but one for which the result is known. In addition, each data set was run with each possible option. Thus roughly 150 approximations have been generated for data sets of typically 12 points.

These tests were performed on IBM 360 model 75 and model 91 computers. Execution time is highly dependent on the data; however the model 91 processed 80 approximation problems of 10 to 12 points each in 2.04 seconds. Since the processing associated with terminating and initiating segments is nonnegligible and since in these constructed test cases the ratio of approximating segments to data points was almost .5, one would expect the program to do slightly better for smoother data.

REFERENCES

1. WILSON, D.G. Piecewise linear approximations of fewest line segments. Proc. AFIPS 1972 SJCC, Vol. 40, AFIPS Press, Montvale, N.J., pp. 187-198.

ALGORITHM

```

SUBROUTINE STL2(X, Y, E, M, U, V, W, K, IP)          STL 10
C PIECEWISE LINEAR APPROXIMATIONS OF FEWEST      STL 20
C LINE SEGMENTS WITHIN GIVEN TOLERANCES.        STL 30
C X,Y,E AND M CONTAIN INPUT DATA.              STL 40
C U,V,K AND POSSIBLY W CONTAIN OUTPUT.          STL 50
C IP IS A PARAMETER DETERMINING THE OPERATION   STL 60
C OF THE PROGRAM.                               STL 70
C X AND Y ARE INPUT DATA ARRAYS OF M ELEMENTS STL 80
C X(I),Y(I) CONTAINS THE ITH DATA POINT.       STL 90
C E MAY BE A SINGLE TOLERANCE OR A TABLE OF   STL 100
C TOLERANCES DEPENDING ON THE VALUE OF IP.     STL 110
C IF E IS AN ARRAY, THEN E(I) IS THE TOLERANCE STL 120
C ASSOCIATED WITH X(I),Y(I) AND E MUST CONTAIN STL 130
C M NONNEGATIVE ELEMENTS.                      STL 140
C U AND V ARE OUTPUT ARRAYS OF K+1 ELEMENTS.   STL 150
C U IS A PARTITION OF THE INTERVAL (X(1),X(N)) STL 160
C WITH U(1)=X(1) AND U(K+1)=X(N).              STL 170
C V(I) IS AN ORDINATE TO BE ASSOCIATED WITH    STL 180
C U(I) IN THE APPROXIMATION. (IF A CONTINUOUS  STL 190
C APPROXIMATION IS REQUESTED, THEN V(I) IS    STL 200
C 'THE' ORDINATE TO BE ASSOCIATED WITH U(I).)  STL 210
C IF A CONTINUOUS APPROXIMATION IS REQUESTED,  STL 220
C THEN W IS NOT USED. IN THIS CASE THE ITH    STL 230
C APPROXIMATING SEGMENT IS THE STRAIGHT LINE  STL 240
C FROM U(I),V(I) TO U(I+1),V(I+1).            STL 250
C IF A CONTINUOUS APPROXIMATION IS NOT        STL 260
C REQUESTED, THEN W IS A K-ELEMENT OUTPUT     STL 270
C ARRAY. IN THIS CASE THE ITH APPROXIMATING   STL 280
C SEGMENT IS THE STRAIGHT LINE FROM           STL 290
C U(I),W(I) TO U(I+1),V(I+1), AND V(1) IS    STL 300
C SET EQUAL TO W(1).                          STL 310
C K IS THE NUMBER OF SEGMENTS IN THE PIECE-   STL 320
C WISE LINEAR APPROXIMATION GENERATED. IN    STL 330
C CASE OF AN ERROR RETURN, K WILL BE SET TO   STL 340
C ZERO.                                        STL 350
C THE CONTROL PARAMETER IP IS THE PRODUCT     STL 360
C OF THREE INDICATORS I1,I2 AND I3.          STL 370
C I1 INDICATES WHETHER OR NOT E IS AN        STL 380
C ARRAY OF TOLERANCES.                       STL 390
C I1 = -1 INDICATES E IS AN ARRAY            STL 400
C I1 = +1 INDICATES E IS A SINGLE NUMBER.    STL 410
C I2 INDICATES WHETHER OR NOT THE           STL 420
C APPROXIMATION IS TO BE RESTRICTED TO      STL 430
C THE 'TOLERANCE BAND' ABOUT THE DATA.     STL 440
C I2 = 1 INDICATES NO BAND RESTRICTION      STL 450
C I2 = 2 INDICATES APPLY THIS RESTRICTION   STL 460
C (THE 'TOLERANCE BAND' IS A PIECEWISE     STL 470
C LINEAR BAND CENTERED AT THE DATA WHOSE  STL 480
C WIDTH IS DETERMINED BY THE TOLERANCES    STL 490
C AT THE DATA POINTS.)                    STL 500
C I3 INDICATES WHETHER OR NOT THE           STL 510
C APPROXIMATION MUST BE CONTINUOUS.        STL 520
C I3 = 1 INDICATES CONTINUITY NOT REQUIRED    STL 530
C I3 = 3 INDICATES CONTINUITY IS REQUIRED    STL 540
C CALL STL2 (X,Y,E,M,X,Y,E,M,IP) WILL NOT   STL 550
C CAUSE PROBLEMS PROVIDED THAT             STL 560
C EITHER A CONTINUOUS APPROXIMATION IS     STL 570
C REQUESTED, OR E IS A SUFFICIENTLY LARGE  STL 580
C ARRAY.                                    STL 590
C THE PROGRAM PERFORMS THE FOLLOWING DATA   STL 600
C CHECKS. ARE THE X-VALUES IN INCREASING    STL 610
C ORDER. ARE THE TOLERANCE(S) NONNEGATIVE.  STL 620
C IS THE NUMBER OF DATA POINTS GREATER    STL 630
C THAN ONE. IF ANY CHECK FAILS, THE PROGRAM  STL 640
C RETURNS WITH K SET EQUAL TO 0. IN THIS    STL 650
C CASE NO FURTHER PROCESSING IS ATTEMPTED.  STL 660
C DIMENSION X(9), Y(9), E(9), U(9), V(9), W(9) STL 670
C N = M                                     STL 680
C ITC = IP                                  STL 690
C J = 1                                     STL 700

```

C ERROR CHECKS	STL 710
IF (N.LE.1) GO TO 400	STL 720
IF (E(1).LT.0.0) GO TO 400	STL 730
DO 10 L=2,N	STL 740
IF (X(L-1).GE.X(L)) GO TO 400	STL 750
IF (ITCH.GE.0) GO TO 10	STL 760
IF (E(L).LT.0.0) GO TO 400	STL 770
10 CONTINUE	STL 780
C INITIALIZATION FOR ENTIRE PROGRAM	STL 790
EPSLN = E(1)	STL 800
SGN = 1.0	STL 810
KEEP = 1	STL 820
I = 1	STL 830
U(1) = X(1)	STL 840
J = 2	STL 850
INIT = 1	STL 860
INDC = 0	STL 870
GO TO 30	STL 880
C INITIALIZATION FOR EACH SEGMENT	STL 890
20 CONTINUE	STL 900
J = J + 1	STL 910
INIT = I	STL 920
INDC = 0	STL 930
IF (IABS(ITCH).LT.3) KEEP = I	STL 940
IF (IABS(IABS(ITCH)-4).NE.2) GO TO 30	STL 950
C RESTRICTED TO TOLERANCE BAND	STL 960
XEYE = U(J-1)	STL 970
YEYE = V(J-1)	STL 980
TEMP1 = EPSLN	STL 990
IF (ITCH.LT.0) TEMP1 = TEMP1 + (SGN*E(I-1)-EPSLN)*(X(I)-U(J-1)	STL 1000
*)/(X(I)-X(I-1))	STL 1010
YINIT = YEYE - TEMP1 - TEMP1	STL 1020
GO TO 40	STL 1030
30 CONTINUE	STL 1040
C NOT RESTRICTED TO TOLERANCE BAND	STL 1050
XEYE = X(I)	STL 1060
YEYE = Y(I) + EPSLN	STL 1070
YINIT = Y(I) - EPSLN	STL 1080
IF (IABS(ITCH).EQ.1 .OR. I.EQ.1) GO TO 40	STL 1090
TEMP1 = EPSLN	STL 1100
IF (ITCH.LT.0) TEMP1 = SGN*E(I+1)	STL 1110
SMIN = (Y(I+1)-YEYE-TEMP1)/(X(I+1)-XEYE)	STL 1120
IF (ITCH.LT.0) TEMP1 = SGN*E(I-1)	STL 1130
SMAX = (YEYE-Y(I-1)+TEMP1)/(XEYE-X(I-1))	STL 1140
IF (KEEP.EQ.I-1) GO TO 50	STL 1150
IT = I - 2	STL 1160
XINIT = XEYE	STL 1170
IPIV = I	STL 1180
IGRAZE = I	STL 1190
I = I + 1	STL 1200
GO TO 150	STL 1210
40 CONTINUE	STL 1220
IF (XEYE.GE.X(I)) I = I + 1	STL 1230
IF (ITCH.LT.0) EPSLN = SGN*E(I)	STL 1240
DX = X(I) - XEYE	STL 1250
SMAX = (Y(I)+EPSLN-YEYE)/DX	STL 1260
SMIN = (Y(I)-EPSLN-YEYE)/DX	STL 1270
50 CONTINUE	STL 1280
XINIT = XEYE	STL 1290
IPIV = I	STL 1300
IGRAZE = I	STL 1310
C DETERMINATION OF INDIVIDUAL SEGMENT	STL 1320
60 CONTINUE	STL 1330
IF (I.EQ.N) GO TO 260	STL 1340
I = I + 1	STL 1350
70 CONTINUE	STL 1360
C TEST FOR NEW *MAX* SLOPE	STL 1370
DX = X(I) - XEYE	STL 1380
IF (ITCH.LT.0) EPSLN = SGN*E(I)	STL 1390
TEMP1 = (Y(I)+EPSLN-YEYE)/DX	STL 1400
TEST = TEMP1 - SMAX	STL 1410
IF (SGN.LE.0.0) TEST = -TEST	STL 1420
IF (TEST) 80, 90, 100	STL 1430
80 CONTINUE	STL 1440
C TEST FOR END OF CANDIDATE SEGMENT	STL 1450

TEST = TEMP1 - SMIN	STL 1460
IF (SGN.LE.0.0) TEST = -TEST	STL 1470
IF (TEST.LT.0.0) GO TO 210	STL 1480
SMAX = TEMP1	STL 1490
90 CONTINUE	STL 1500
C TEST FOR NEW *MIN* SLOPE	STL 1510
IPIV = I	STL 1520
100 CONTINUE	STL 1530
TEMP2 = (Y(I)-EPSLN-YEYE)/DX	STL 1540
TEST = TEMP2 - SMAX	STL 1550
IF (SGN.LE.0.0) TEST = -TEST	STL 1560
IF (TEST) 110, 120, 140	STL 1570
110 CONTINUE	STL 1580
TEST = SMIN - TEMP2	STL 1590
IF (SGN.LE.0.0) TEST = -TEST	STL 1600
IF (TEST) 120, 130, 60	STL 1610
120 CONTINUE	STL 1620
SMIN = TEMP2	STL 1630
130 CONTINUE	STL 1640
IGRAZE = I	STL 1650
GO TO 60	STL 1660
C CHECK FOR PIVOT AT NEW EYE POINT	STL 1670
140 CONTINUE	STL 1680
IF (XEYE.EQ.X(IPIV)) GO TO 220	STL 1690
IF (ITCH.LT.0) EPSLN = SGN*E(IPIV)	STL 1700
INDC = 1	STL 1710
SVX = XEYE	STL 1720
SVY = YEYE	STL 1730
SVMN = SMIN	STL 1740
SVMX = SMAX	STL 1750
XEYE = X(IPIV)	STL 1760
YEYE = Y(IPIV) + EPSLN	STL 1770
SMIN = SMAX	STL 1780
SMAX = (YINIT-YEYE)/(XINIT-XEYE)	STL 1790
IF (KEEP.GE.IPIV) GO TO 170	STL 1800
IT = IPIV - 1	STL 1810
150 CONTINUE	STL 1820
TEMP2 = YEYE + EPSLN	STL 1830
DO 160 L=KEEP,IT	STL 1840
IF (ITCH.LT.0) TEMP2 = YEYE + SGN*E(L)	STL 1850
TEMP1 = (Y(L)-TEMP2)/(X(L)-XEYE)	STL 1860
TEST = TEMP1 - SMAX	STL 1870
IF (SGN.LE.0.0) TEST = -TEST	STL 1880
IF (TEST.LT.0.0) SMAX = TEMP1	STL 1890
160 CONTINUE	STL 1900
170 CONTINUE	STL 1910
IF (IPIV.GE.I-1) GO TO 70	STL 1920
IT = I - 2	STL 1930
TEMP2 = YEYE - EPSLN	STL 1940
IDIOT = IPIV	STL 1950
DO 200 L=IDIOT,IT	STL 1960
DX = X(L+1) - XEYE	STL 1970
IF (ITCH.LT.0) TEMP2 = YEYE - SGN*E(L+1)	STL 1980
TEMP1 = (Y(L+1)-TEMP2)/DX	STL 1990
TEST = TEMP1 - SMAX	STL 2000
IF (SGN.LE.0.0) TEST = -TEST	STL 2010
IF (TEST) 180, 190, 200	STL 2020
180 CONTINUE	STL 2030
SMAX = TEMP1	STL 2040
190 CONTINUE	STL 2050
IPIV = L + 1	STL 2060
200 CONTINUE	STL 2070
GO TO 70	STL 2080
C END OF CURRENT SEGMENT	STL 2090
210 CONTINUE	STL 2100
TEMP2 = SMIN	STL 2110
IF (I.EQ.N) GO TO 240	STL 2120
KEEP = IGRAZE	STL 2130
GO TO 250	STL 2140
220 CONTINUE	STL 2150
TEMP2 = SMAX	STL 2160
IF (I.EQ.N) GO TO 230	STL 2170
SGN = -SGN	STL 2180
EPSLN = -EPSLN	STL 2190
KEEP = IPIV	STL 2200
GO TO 250	STL 2210

230	CONTINUE	STL 2220
	IF (INDC.EQ.0 .OR. XEYE.NE.X(N-1)) GO TO 240	STL 2230
	XEYE = SVX	STL 2240
	YEYE = SVY	STL 2250
	SMIN = SVMN	STL 2260
	SMAX = SVMX	STL 2270
240	CONTINUE	STL 2280
	U(J) = X(N-1)	STL 2290
	YINIT = Y(N-1)	STL 2300
	GO TO 270	STL 2310
250	CONTINUE	STL 2320
	IF (IABS(IABS(ITCH)-4).NE.2) GO TO 300	STL 2330
C	DETERMINE KNOT ON EDGE OF TOLERANCE BAND	STL 2340
	TEMP1 = 0.0	STL 2350
	IF (ITCH.LT.0) TEMP1 = EPSLN - SGN*E(I-1)	STL 2360
	TEMP1 = (Y(I)-Y(I-1)+TEMP1)/(X(I)-X(I-1))	STL 2370
	U(J) = (Y(I)+EPSLN-YEYE-TEMP1*X(I)+TEMP2*XEYE)/(TEMP2-TEMP1)	STL 2380
	GO TO 310	STL 2390
260	CONTINUE	STL 2400
	U(J) = X(N)	STL 2410
	YINIT = Y(N)	STL 2420
270	CONTINUE	STL 2430
C	CONTINUITY CHECK FOR LAST SEGMENT	STL 2440
	IF (IABS(ITCH).GE.3 .OR. INIT.EQ.1) GO TO 290	STL 2450
	IT = INIT - 1	STL 2460
	SVMX = SMAX + SGN	STL 2470
	TEMP2 = YEYE + EPSLN	STL 2480
	DO 280 L=KP,IT	STL 2490
	IF (ITCH.LT.0) TEMP2 = YEYE + SGN*E(L)	STL 2500
	TEMP1 = (Y(L)-TEMP2)/(X(L)-XEYE)	STL 2510
	TEST = TEMP1 - SVMX	STL 2520
	IF (SGN.LE.0.0) TEST = -TEST	STL 2530
	IF (TEST.LT.0.0) SVMX = TEMP1	STL 2540
280	CONTINUE	STL 2550
	IF (ABS(SVMX-SMAX+SVMX-SMIN).LE.ABS(SMAX-SMIN)) SMAX = SVMX	STL 2560
290	CONTINUE	STL 2570
C	NEARNESS CHECK FOR LAST SEGMENT	STL 2580
	TEMP2 = SMAX	STL 2590
	TEMP1 = YEYE + SMAX*(U(J)-XEYE)	STL 2600
	TEST = YINIT - TEMP1	STL 2610
	IF (SGN.LT.0.0) TEST = -TEST	STL 2620
	IF (TEST.GT.0.0) GO TO 310	STL 2630
	TEMP2 = SMIN	STL 2640
	TEMP1 = YEYE + SMIN*(U(J)-XEYE)	STL 2650
	TEST = YINIT - TEMP1	STL 2660
	IF (SGN.LT.0.0) TEST = -TEST	STL 2670
	IF (TEST.LT.0.0) GO TO 310	STL 2680
	TEMP2 = (YINIT-YEYE)/(U(J)-XEYE)	STL 2690
	V(J) = YINIT	STL 2700
	GO TO 320	STL 2710
300	CONTINUE	STL 2720
	IF (IABS(ITCH).GE.3) GO TO 330	STL 2730
	U(J) = 0.5*(X(I)+X(I-1))	STL 2740
310	CONTINUE	STL 2750
	V(J) = YEYE + TEMP2*(U(J)-XEYE)	STL 2760
320	CONTINUE	STL 2770
	IF (XEYE.NE.XINIT) GO TO 330	STL 2780
	IF (IABS(ITCH).EQ.2) GO TO 360	STL 2790
	IF (IABS(ITCH).NE.6) GO TO 330	STL 2800
	IF (J.LE.2) GO TO 380	STL 2810
	GO TO 390	STL 2820
330	CONTINUE	STL 2830
C	RECOMPUTATION OF KNOT FOR CONTINUITY	STL 2840
	IF (J.LE.2) GO TO 370	STL 2850
	IF (SLOPE.EQ.TEMP2) GO TO 360	STL 2860
	YINIT = V(J-2)	STL 2870
	IF (IABS(ITCH).LT.3) YINIT = W(J-2)	STL 2880
	TEMP1 = (XEYE*TEMP2-U(J-2)*SLOPE+YINIT-YEYE)/(TEMP2-SLOPE)	STL 2890
	IF (IABS(ITCH).GE.3) GO TO 350	STL 2900
	IF (TEMP1.GT.XINIT) GO TO 360	STL 2910
	TEST = ABS(EPSLN)	STL 2920
	IDIOT = INIT - KP	STL 2930
	DO 340 L=1,IDIOT	STL 2940
	IT = INIT - L	STL 2950
	IF (TEMP1.GE.X(IT)) GO TO 350	STL 2960
	DX = Y(IT) - YEYE - TEMP2*(X(IT)-XEYE)	STL 2970

IF (ITCH.LT.0) TEST = E(IT)	STL 2980
IF (ABS(DX).GT.TEST) GO TO 360	STL 2990
340 CONTINUE	STL 3000
350 CONTINUE	STL 3010
U(J-1) = TEMP1	STL 3020
V(J-1) = YEYE + TEMP2*(U(J-1)-XEYE)	STL 3030
IF (IABS(ITCH).LT.3) W(J-1) = V(J-1)	STL 3040
GO TO 390	STL 3050
360 CONTINUE	STL 3060
W(J-1) = YEYE + TEMP2*(U(J-1)-XEYE)	STL 3070
GO TO 390	STL 3080
370 CONTINUE	STL 3090
IF (IABS(ITCH).LT.3) GO TO 360	STL 3100
380 CONTINUE	STL 3110
V(1) = YEYE + TEMP2*(U(1)-XEYE)	STL 3120
390 CONTINUE	STL 3130
SLOPE = TEMP2	STL 3140
KP = KEEP	STL 3150
IF (I.LT.N) GO TO 20	STL 3160
IF (X(N).EQ.U(J)) GO TO 400	STL 3170
IF (IABS(ITCH).LT.3) W(J) = V(J)	STL 3180
J = J + 1	STL 3190
U(J) = X(N)	STL 3200
V(J) = Y(N)	STL 3210
400 CONTINUE	STL 3220
IF (J.GE.2 .AND. IABS(ITCH).LT.3) V(1) = W(1)	STL 3230
K = J - 1	STL 3240
RETURN	STL 3250
END	STL 3260

ALGORITHM 511
 CDC 6600 Subroutines IBESS and JBESS
 for Bessel Functions
 $I_\nu(x)$ and $J_\nu(x)$, $x \geq 0$, $\nu \geq 0$ [S18]

D. E. AMOS, S. L. DANIEL, and M. K. WESTON
 Sandia Laboratories

Key Words and Phrases: Bessel functions of the first kind, modified Bessel function, Airy functions, uniform asymptotic expansion
 CR Categories: 5.12
 Language: Fortran

DESCRIPTION

This algorithm is a complement to [1], where the theoretical background and development are described.

REFERENCES

1. AMOS, D.E., DANIEL, S.L., AND WESTON, M.K. CDC 6600 Subroutines IBESS and JBESS for Bessel Functions $I_\nu(x)$ and $J_\nu(x)$, $x \geq 0$, $\nu \geq 0$. *ACM Trans. Math. Software* 3, 1 (March 1977), 76-92.

ALGORITHM

	FUNCTION GAMLN(X)	GAM	10
C		GAM	20
C	A CDC 6600 SUBROUTINE	GAM	30
C		GAM	40
C	AUTHORS	GAM	50
C	D.E. AMOS AND S.L. DANIEL	GAM	60
C	ALBUQUERQUE, NEW MEXICO, 87115	GAM	70
C	JANUARY, 1975	GAM	80
C		GAM	90
C	REFERENCES	GAM	100
C	ABRAMOWITZ, M. AND STEGUN, I.A. HANDBOOK OF MATHEMATICAL	GAM	110
C	FUNCTIONS. NBS APPLIED MATHEMATICS SERIES 55, U.S. GOVERNMENT	GAM	120
C	PRINTING OFFICE, WASHINGTON, D.C., CHAPTER 6.	GAM	130
C		GAM	140
C	AMOS, D.E., DANIEL, S.L. AND WESTON, M.K. CDC 6600	GAM	150
C	SUBROUTINES IBESS AND JBESS FOR BESSEL FUNCTIONS	GAM	160
C	I/SUB(NU)/(X) AND J/SUB(NU)/(X), X.GE.0, NU.GE.0.	GAM	170
C	ACM TRANS. MATH. SOFTWARE, 1977.	GAM	180
C		GAM	190
C	HART, J.F., ET. AL. COMPUTER APPROXIMATIONS, WILEY, NEW YORK.	GAM	200
C	PP. 130-136, 1968.	GAM	210
C		GAM	220
C	ABSTRACT	GAM	230

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130	GAMLN = G(8)	GAM	1000
	RETURN	GAM	1010
140	PRINT 99999, X	GAM	1020
	STOP	GAM	1030
99999	FORMAT (50H ARGUMENT FOR GAMLN IS LESS THAN OR EQUAL TO ZERO,,	GAM	1040
	* 3H X=,E25.14)	GAM	1050
	END	GAM	1060
	SUBROUTINE IBESS(KODE, ALPHA, N, X, Y)	IBE	10
C		IBE	20
C	A CDC 6600 SUBROUTINE	IBE	30
C		IBE	40
C	AUTHORS	IBE	50
C	D.E. AMOS AND S.L. DANIEL	IBE	60
C	SANDIA LABORATORIES	IBE	70
C	JANUARY, 1975	IBE	80
C		IBE	90
C	REFERENCES	IBE	100
C	ABRAMOWITZ, M. AND STEGUN, I.A. HANDBOOK OF MATHEMATICAL	IBE	110
C	FUNCTIONS. NBS APPLIED MATHEMATICS SERIES 55, U.S. GOVERNMENT	IBE	120
C	PRINTING OFFICE, WASHINGTON, D.C., CHAPTERS 9 AND 10.	IBE	130
C		IBE	140
C	AMOS, D.E., DANIEL, S.L. AND WESTON, M.K. CDC 6600	IBE	150
C	SUBROUTINES IBESS AND JBESS FOR BESSEL FUNCTIONS	IBE	160
C	I/SUB(NU)/(X) AND J/SUB(NU)/(X), X.GE.0, NU.GE.0.	IBE	170
C	ACM TRANS. MATH. SOFTWARE, 1977.	IBE	180
C		IBE	190
C	OLVER, F.W.J. TABLES FOR BESSEL FUNCTIONS OF MODERATE OR	IBE	200
C	LARGE ORDERS. NPL MATHEMATICAL TABLES, VOL 6. HER MAJESTY-S	IBE	210
C	STATIONERY OFFICE, LONDON, 1962.	IBE	220
C		IBE	230
C	OLVER, F.W.J. THE ASYMPTOTIC EXPANSION OF BESSEL FUNCTIONS OF	IBE	240
C	LARGE ORDER. PHIL. TRANS. A, 247, PP. 328-368, 1954.	IBE	250
C		IBE	260
C	ABSTRACT	IBE	270
C	IBESS COMPUTES AN N MEMBER SEQUENCE OF I BESSEL FUNCTIONS	IBE	280
C	I/SUB(ALPHA+K-1)/(X), K=1,...,N OR SCALED BESSEL FUNCTIONS	IBE	290
C	EXP(-X)*I/SUB(ALPHA*K-1)/(X), K=1,...,N FOR NON-NEGATIVE ALPHA	IBE	300
C	AND X. A COMBINATION OF THE POWER SERIES, THE ASYMPTOTIC	IBE	310
C	EXPANSION FOR X TO INFINITY, AND THE UNIFORM ASYMPTOTIC	IBE	320
C	EXPANSION FOR NU TO INFINITY ARE APPLIED OVER SUBDIVISIONS OF	IBE	330
C	THE (NU,X) PLANE. FOR VALUES NOT COVERED BY ONE OF THESE	IBE	340
C	FORMULAE, THE ORDER IS INCREMENTED BY AN INTEGER SO THAT ONE	IBE	350
C	OF THESE FORMULAE APPLY. BACKWARD RECURSION IS USED TO REDUCE	IBE	360
C	ORDERS BY INTEGER VALUES. THE ASYMPTOTIC EXPANSION FOR X TO	IBE	370
C	INFINITY IS USED ONLY WHEN THE ENTIRE SEQUENCE (SPECIFICALLY	IBE	380
C	THE LAST MEMBER) LIES WITHIN THE REGION COVERED BY THE	IBE	390
C	EXPANSION. LEADING TERMS OF THESE EXPANSIONS ARE USED TO TEST	IBE	400
C	FOR OVER OR UNDERFLOW WHERE APPROPRIATE. IF A SEQUENCE IS	IBE	410
C	REQUESTED AND THE LAST MEMBER WOULD UNDERFLOW, THE RESULT IS	IBE	420
C	SET TO ZERO AND THE NEXT LOWER ORDER TRIED, ETC., UNTIL A	IBE	430
C	MEMBER COMES ON SCALE OR ALL ARE SET TO ZERO. AN OVERFLOW	IBE	440
C	CANNOT OCCUR WITH SCALING. IBESS CALLS FUNCTION GAMLN.	IBE	450
C		IBE	460
C	DESCRIPTION OF ARGUMENTS	IBE	470
C		IBE	480
C	INPUT	IBE	490
C	KODE - A PARAMETER TO INDICATE THE SCALING OPTION	IBE	500
C	KODE=1 RETURNS	IBE	510
C	Y(K)= I/SUB(ALPHA+K-1)/(X),	IBE	520
C	K=1,...,N	IBE	530
C	KODE=2 RETURNS	IBE	540
C	Y(K)=EXP(-X)*I/SUB(ALPHA+K-1)/(X),	IBE	550
C	K=1,...,N	IBE	560
C	ALPHA - ORDER OF FIRST MEMBER OF THE SEQUENCE, ALPHA.GE.0	IBE	570
C	N - NUMBER OF MEMBERS IN THE SEQUENCE, N.GE.1	IBE	580
C	X - X.GE.0	IBE	590
C		IBE	600
C	OUTPUT	IBE	610
C	Y - A VECTOR WHOSE FIRST N COMPONENTS CONTAIN	IBE	620
C	VALUES FOR I/SUB(ALPHA+K-1)/(X) OR SCALED	IBE	630
C	VALUES FOR EXP(-X)*I/SUB(ALPHA+K-1)/(X),	IBE	640
C	K=1,...,N DEPENDING ON KODE	IBE	650
C		IBE	660
C	ERROR CONDITIONS	IBE	670

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C      IMPROPER INPUT ARGUMENTS - A FATAL ERROR      IBE 680
C      OVERFLOW WITH KODE=1 - A FATAL ERROR          IBE 690
C      UNDERFLOW - A NON-FATAL ERROR                 IBE 700
C                                                     IBE 710
      DOUBLE PRECISION DX, TRX, DTM, DFN             IBE 720
      DIMENSION Y(1), TEMP(3)                        IBE 730
      DIMENSION C(11,10)                             IBE 740
      DIMENSION C1(11,8), C2(11,2)                  IBE 750
      EQUIVALENCE (C(1,1),C1(1,1))                  IBE 760
      EQUIVALENCE (C(1,9),C2(1,1))                  IBE 770
      DATA ELIM,TOL / 667., 1.E-15 / IBE 780
      DATA RTP1,RTTPI / 1.59154943091895E-01, 3.98942280401433E-01/ IBE 790
      DATA CE / 3.45387763900000E+01/ IBE 800
      DATA INLIM / 80 / IBE 810
      DATA C1 / -2.08333333333333E-01, 1.25000000000000E-01, IBE 820
1      9(0.) , 3.34201388888888E-01, -4.01041666666667E-01, IBE 830
2      7.03125000000000E-02, 8(0.) , -1.02581259645062E+00, IBE 840
3      1.84646267361111E+00, -8.91210937500000E-01, 7.32421875000000E-02, IBE 850
4      7(0.) , 4.66958442342625E+00, -1.12070026162230E+01, IBE 860
5      8.78912353515625E+00, -2.36408691406250E+00, 1.12152099609375E-01, IBE 870
6      6(0.) , -2.82120725582002E+01, 8.46362176746007E+01, IBE 880
7-9.18182415432400E+01, 4.25349987453885E+01, -7.36879435947963E+00, IBE 890
8      2.27108001708984E-01, 5(0.) , 2.12570130039217E+02, IBE 900
9-7.65252468141182E+02, 1.05999045252800E+03, -6.99579627376133E+02, IBE 910
A      2.18190511744212E+02, -2.64914304869516E+01, 5.72501420974731E-01, IBE 920
B      4(0.) , -1.91945766231841E+03, 8.06172218173731E+03, IBE 930
C-1.35865500064341E+04, 1.16553933368645E+04, -5.30564697861340E+03, IBE 940
D      1.20090291321635E+03, -1.08090919788395E+02, 1.72772750258446E+00, IBE 950
E      3(0.) , 2.02042913309661E+04, -9.69805983886375E+04, IBE 960
F      1.92547001232532E+05, -2.03400177280416E+05, 1.22200464983017E+05, IBE 970
G-4.11926549688976E+04, 7.10951430248936E+03, -4.93915304773088E+02, IBE 980
H      6.07404200127348E+00, 2(0.) / IBE 990
      DATA C2 / -2.42919187900551E+05, 1.31176361466298E+06, IBE 1000
1-2.99801591853811E+06, 3.76327129765640E+06, -2.81356322658653E+06, IBE 1010
2      1.26836527332162E+06, -3.31645172484564E+05, 4.52187689813627E+04, IBE 1020
3-2.49983048181121E+03, 2.43805296995561E+01, 1(0.) , IBE 1030
4      3.28446985307204E+06, -1.97068191184322E+07, 5.09526024926646E+07, IBE 1040
5-7.41051482115327E+07, 6.63445122747290E+07, -3.75671766607634E+07, IBE 1050
6      1.32887671664218E+07, -2.78561812808645E+06, 3.08186404612662E+05, IBE 1060
7-1.38860897537170E+04, 1.10017140269247E+02/ IBE 1070
C TEST INPUT ARGUMENTS IBE 1080
      KT = 1 IBE 1090
C TEST INPUT ARGUMENTS IBE 1100
      IF (N-1) 580, 10, 20 IBE 1110
10     KT = 2 IBE 1120
20     NN = N IBE 1130
      IF (KODE.LT.1 .OR. KODE.GT.2) GO TO 560 IBE 1140
      IF (X) 590, 30, 80 IBE 1150
30     IF (ALPHA) 570, 40, 50 IBE 1160
40     Y(1) = 1. IBE 1170
      IF (N.EQ.1) RETURN IBE 1180
      I1 = 2 IBE 1190
      GO TO 60 IBE 1200
50     I1 = 1 IBE 1210
60     DO 70 I=I1,N IBE 1220
          Y(I) = 0. IBE 1230
70     CONTINUE IBE 1240
      RETURN IBE 1250
80     CONTINUE IBE 1260
      IF (ALPHA.LT.0.) GO TO 570 IBE 1270
      DFN = DBLE(FLOAT(N)) + DBLE(ALPHA) - 1.D+0 IBE 1280
      FNU = DFN IBE 1290
      IN = 0 IBE 1300
      XO2 = X*.5 IBE 1310
      SXO2 = XO2*XO2 IBE 1320
      ETX = KODE - 1 IBE 1330
      SX = ETX*X IBE 1340
C DECISION TREE FOR REGION WHERE SERIES, ASYMPTOTIC EXPANSION FOR X IBE 1350
C TO INFINITY AND ASYMPTOTIC EXPANSION FOR NU TO INFINITY ARE IBE 1360
C APPLIED. IBE 1370
      IF (SXO2.LE.(FNU+1.)) GO TO 90 IBE 1380
      IF (X.LE.12.) GO TO 110 IBE 1390
      FN = 0.55*FNU*FNU IBE 1400
      FN = AMAX1(17.,FN) IBE 1410
      IF (X.GE.FN) GO TO 430 IBE 1420
      NS = AMAX1(36.-FNU,0.) IBE 1430

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DFN = DFN + DBLE(FLOAT(NS))	IBE 1440
FN = DFN	IBE 1450
IS = KT	IBE 1460
KM = N - 1 + NS	IBE 1470
IF (KM.GT.0) IS = 3	IBE 1480
GO TO 120	IBE 1490
90 FN = FNU	IBE 1500
FNP1 = FN + 1.	IBE 1510
XO2L = ALOG(XO2)	IBE 1520
IS = KT	IBE 1530
IF (X.LE.0.5) GO TO 230	IBE 1540
NS = 0	IBE 1550
100 DFN = DFN + DBLE(FLOAT(NS))	IBE 1560
FN = DFN	IBE 1570
FNP1 = FN + 1.	IBE 1580
IS = KT	IBE 1590
IF (N-1+NS.GT.0) IS = 3	IBE 1600
GO TO 230	IBE 1610
110 XO2L = ALOG(XO2)	IBE 1620
NS = SXO2 - FNU	IBE 1630
GO TO 100	IBE 1640
120 CONTINUE	IBE 1650
C OVERFLOW TEST ON UNIFORM ASYMPTOTIC EXPANSION	IBE 1660
IF (KODE.EQ.2) GO TO 130	IBE 1670
IF (ALPHA.LT.1.) GO TO 150	IBE 1680
Z = X/ALPHA	IBE 1690
RA = SQRT(1.+Z*Z)	IBE 1700
GLN = ALOG((1.+RA)/Z)	IBE 1710
T = RA*(1.-ETX) + ETX/(Z+RA)	IBE 1720
ARG = ALPHA*(T-GLN)	IBE 1730
IF (ARG.GT.ELIM) GO TO 600	IBE 1740
IF (KM.EQ.0) GO TO 140	IBE 1750
130 CONTINUE	IBE 1760
C UNDERFLOW TEST ON UNIFORM ASYMPTOTIC EXPANSION	IBE 1770
Z = X/FN	IBE 1780
RA = SQRT(1.+Z*Z)	IBE 1790
GLN = ALOG((1.+RA)/Z)	IBE 1800
T = RA*(1.-ETX) + ETX/(Z+RA)	IBE 1810
ARG = FN*(T-GLN)	IBE 1820
140 IF (ARG.LT.-ELIM) GO TO 280	IBE 1830
GO TO 190	IBE 1840
150 IF (X.GT.ELIM) GO TO 600	IBE 1850
GO TO 130	IBE 1860
C UNIFORM ASYMPTOTIC EXPANSION FOR NU TO INFINITY	IBE 1870
160 IF (KM.NE.0) GO TO 170	IBE 1880
Y(1) = TEMP(3)	IBE 1890
RETURN	IBE 1900
170 TEMP(1) = TEMP(3)	IBE 1910
IN = NS	IBE 1920
KT = 1	IBE 1930
180 CONTINUE	IBE 1940
IS = 2	IBE 1950
DFN = DFN - 1.D+0	IBE 1960
FN = DFN	IBE 1970
Z = X/FN	IBE 1980
RA = SQRT(1.+Z*Z)	IBE 1990
GLN = ALOG((1.+RA)/Z)	IBE 2000
T = RA*(1.-ETX) + ETX/(Z+RA)	IBE 2010
ARG = FN*(T-GLN)	IBE 2020
190 COEF = EXP(ARG)	IBE 2030
T = 1./RA	IBE 2040
T2 = T*T	IBE 2050
T = T/FN	IBE 2060
S2 = 1.	IBE 2070
AP = 1.	IBE 2080
DO 210 K=1,10	IBE 2090
KP1 = K + 1	IBE 2100
S1 = C(1,K)	IBE 2110
DO 200 J=2,KP1	IBE 2120
S1 = S1*T2 + C(J,K)	IBE 2130
200 CONTINUE	IBE 2140
AP = AP*T	IBE 2150
AK = AP*S1	IBE 2160
S2 = S2 + AK	IBE 2170
IF (ABS(AK).LT.TOL) GO TO 220	IBE 2180
210 CONTINUE	IBE 2190

220	CONTINUE	IBE 2200
	TEMP(IS) = SQRT(T*RTPI)*COEF*S2	IBE 2210
	GO TO (180, 350, 500), IS	IBE 2220
C	SERIES FOR (X/2)**2.LE.NU+1	IBE 2230
230	CONTINUE	IBE 2240
	GLN = GMLN(FNP1)	IBE 2250
	ARG = FN*X02L - GLN - SX	IBE 2260
	IF (ARG.LT.-ELIM) GO TO 300	IBE 2270
	EARG = EXP(ARG)	IBE 2280
240	CONTINUE	IBE 2290
	S = 1.	IBE 2300
	AK = 3.	IBE 2310
	T2 = 1.	IBE 2320
	T = 1.	IBE 2330
	S1 = FN	IBE 2340
	DO 250 K=1,17	IBE 2350
	S2 = T2 + S1	IBE 2360
	T = T*SX02/S2	IBE 2370
	S = S + T	IBE 2380
	IF (ABS(T).LT.TOL) GO TO 260	IBE 2390
	T2 = T2 + AK	IBE 2400
	AK = AK + 2.	IBE 2410
	S1 = S1 + FN	IBE 2420
250	CONTINUE	IBE 2430
260	CONTINUE	IBE 2440
	TEMP(IS) = S*EARG	IBE 2450
	GO TO (270, 350, 490), IS	IBE 2460
270	EARG = EARG*FN/X02	IBE 2470
	DFN = DFN - 1.D+0	IBE 2480
	FN = DFN	IBE 2490
	IS = 2	IBE 2500
	GO TO 240	IBE 2510
C	SET UNDERFLOW VALUE AND UPDATE PARAMETERS	IBE 2520
280	Y(NN) = 0.	IBE 2530
	NN = NN - 1	IBE 2540
	DFN = DFN - 1.D+0	IBE 2550
	FN = DFN	IBE 2560
	IF (NN-1) 340, 290, 130	IBE 2570
290	KT = 2	IBE 2580
	IS = 2	IBE 2590
	GO TO 130	IBE 2600
300	Y(NN) = 0.	IBE 2610
	NN = NN - 1	IBE 2620
	FNP1 = FN	IBE 2630
	DFN = DFN - 1.D+0	IBE 2640
	FN = DFN	IBE 2650
	IF (NN-1) 340, 310, 320	IBE 2660
310	KT = 2	IBE 2670
	IS = 2	IBE 2680
320	IF (SX02.LE.FNP1) GO TO 330	IBE 2690
	GO TO 130	IBE 2700
330	ARG = ARG - X02L + ALOG(FNP1)	IBE 2710
	IF (ARG.LT.-ELIM) GO TO 300	IBE 2720
	GO TO 230	IBE 2730
340	NZ = N - NN	IBE 2740
	PRINT 99994, NZ, KODE, ALPHA, N, X	IBE 2750
	RETURN	IBE 2760
C	BACKWARD RECURSION SECTION	IBE 2770
350	CONTINUE	IBE 2780
	NZ = N - NN	IBE 2790
	IF (NZ.NE.0) PRINT 99994, NZ, KODE, ALPHA, N, X	IBE 2800
360	GO TO (370, 420), KT	IBE 2810
370	CONTINUE	IBE 2820
	S1 = TEMP(1)	IBE 2830
	S2 = TEMP(2)	IBE 2840
	DX = X	IBE 2850
	TRX = 2.D+0/DX	IBE 2860
	DTM = DFN*TRX	IBE 2870
	TM = DTM	IBE 2880
	IF (IN.EQ.0) GO TO 390	IBE 2890
C	BACKWARD RECUR TO INDEX ALPHA+NN-1	IBE 2900
	DO 380 I=1,IN	IBE 2910
	S = S2	IBE 2920
	S2 = TM*S2 + S1	IBE 2930
	S1 = S	IBE 2940
	DTM = DTM - TRX	IBE 2950

TM = DTM	IBE 2960
380 CONTINUE	IBE 2970
Y(NN) = S1	IBE 2980
IF (NN.EQ.1) RETURN	IBE 2990
Y(NN-1) = S2	IBE 3000
IF (NN.EQ.2) RETURN	IBE 3010
GO TO 400	IBE 3020
390 CONTINUE	IBE 3030
C BACKWARD RECUR FROM INDEX ALPHA+NN-1 TO ALPHA	IBE 3040
Y(NN) = S1	IBE 3050
Y(NN-1) = S2	IBE 3060
IF (NN.EQ.2) RETURN	IBE 3070
400 K = NN + 1	IBE 3080
DO 410 I=3,NN	IBE 3090
K = K - 1	IBE 3100
Y(K-2) = TM*Y(K-1) + Y(K)	IBE 3110
DTM = DTM - TRX	IBE 3120
TM = DTM	IBE 3130
410 CONTINUE	IBE 3140
RETURN	IBE 3150
420 Y(1) = TEMP(2)	IBE 3160
RETURN	IBE 3170
C ASYMPTOTIC EXPANSION FOR X TO INFINITY	IBE 3180
430 CONTINUE	IBE 3190
EARG = RTPI/SQRT(X)	IBE 3200
IF (KODE.EQ.2) GO TO 440	IBE 3210
IF (X.GT.ELIM) GO TO 600	IBE 3220
EARG = EARG*EXP(X)	IBE 3230
440 ETX = 8.*X	IBE 3240
IS = KT	IBE 3250
IN = 0	IBE 3260
450 DX = DFN + DFN	IBE 3270
DTM = DX*DX	IBE 3280
S1 = ETX	IBE 3290
TRX = S1	IBE 3300
DX = -(DTM-1.D+0)/TRX	IBE 3310
T = DX	IBE 3320
TRX = 1.D+0 + DX	IBE 3330
S = TRX	IBE 3340
S2 = 1.	IBE 3350
AK = 8.	IBE 3360
DO 460 K=1,25	IBE 3370
S1 = S1 + ETX	IBE 3380
S2 = S2 + AK	IBE 3390
DX = S2	IBE 3400
TRX = DTM - DX	IBE 3410
AP = TRX	IBE 3420
T = -T*AP/S1	IBE 3430
S = S + T	IBE 3440
IF (ABS(T).LE.TOL) GO TO 470	IBE 3450
AK = AK + 8.	IBE 3460
460 CONTINUE	IBE 3470
470 TEMP(IS) = S*EARG	IBE 3480
GO TO (480, 360), IS	IBE 3490
480 IS = 2	IBE 3500
DFN = DFN - 1.D+0	IBE 3510
GO TO 450	IBE 3520
C BACKWARD RECURSION WITH NORMALIZATION BY	IBE 3530
C ASYMPTOTIC EXPANSION FOR NU TO INFINITY OR POWER SERIES.	IBE 3540
490 CONTINUE	IBE 3550
C COMPUTATION OF LAST ORDER FOR SERIES NORMALIZATION	IBE 3560
KM = AMAX1(3.-FN,0.)	IBE 3570
TFN = FN + FLOAT(KM)	IBE 3580
TA = (GLN+TFN-0.9189385332-0.0833333333/TFN)/(TFN+0.5)	IBE 3590
TA = XO2L - TA	IBE 3600
TB = -(1.-1./TFN)/TFN	IBE 3610
IN = CE/(-TA+SQRT(TA*TA-CE*TB)) + 1.5	IBE 3620
IN = IN + KM	IBE 3630
GO TO 510	IBE 3640
500 CONTINUE	IBE 3650
C COMPUTATION OF LAST ORDER FOR ASYMPTOTIC EXPANSION NORMALIZATION	IBE 3660
IN = CE/(GLN+SQRT(GLN*GLN+T*CE)) + 1.5	IBE 3670
IF (IN.GT.INLIM) GO TO 160	IBE 3680
510 DX = FLOAT(IN)	IBE 3690
DTM = DFN + DX	IBE 3700
DX = X	IBE 3710

TRX = 2.D+0/DX	IBE 3720
DTM = DTM*TRX	IBE 3730
TM = DTM	IBE 3740
TA = 0.	IBE 3750
TB = TOL	IBE 3760
KK = 1	IBE 3770
520 CONTINUE	IBE 3780
C BACKWARD RECUR UNINDEXED	IBE 3790
DO 530 I=1,IN	IBE 3800
S = TB	IBE 3810
TB = TM*TB + TA	IBE 3820
TA = S	IBE 3830
DTM = DTM - TRX	IBE 3840
TM = DTM	IBE 3850
530 CONTINUE	IBE 3860
C NORMALIZATION	IBE 3870
IF (KK.NE.1) GO TO 540	IBE 3880
TA = (TA/TB)*TEMP(3)	IBE 3890
TB = TEMP(3)	IBE 3900
KK = 2	IBE 3910
IN = NS	IBE 3920
IF (NS.NE.0) GO TO 520	IBE 3930
540 Y(NN) = TB	IBE 3940
NZ = N - NN	IBE 3950
IF (NZ.NE.0) PRINT 99994, NZ, KODE, ALPHA, N, X	IBE 3960
IF (NN.EQ.1) RETURN	IBE 3970
TB = TM*TB + TA	IBE 3980
K = NN - 1	IBE 3990
Y(K) = TB	IBE 4000
IF (NN.EQ.2) RETURN	IBE 4010
DTM = DTM - TRX	IBE 4020
TM = DTM	IBE 4030
KM = K - 1	IBE 4040
C BACKWARD RECUR INDEXED	IBE 4050
DO 550 I=1,KM	IBE 4060
Y(K-1) = TM*Y(K) + Y(K+1)	IBE 4070
DTM = DTM - TRX	IBE 4080
TM = DTM	IBE 4090
K = K - 1	IBE 4100
550 CONTINUE	IBE 4110
RETURN	IBE 4120
560 PRINT 99999, KODE, ALPHA, N, X	IBE 4130
STOP	IBE 4140
570 PRINT 99998, KODE, ALPHA, N, X	IBE 4150
STOP	IBE 4160
580 PRINT 99997, KODE, ALPHA, N, X	IBE 4170
STOP	IBE 4180
590 PRINT 99996, KODE, ALPHA, N, X	IBE 4190
STOP	IBE 4200
600 PRINT 99995, KODE, ALPHA, N, X	IBE 4210
STOP	IBE 4220
99999 FORMAT (51H0IBESS CALLED WITH SCALING OPTION, KODE, NOT 1 OR 2	IBE 4230
* /6H KODE=,I2,7H ALPHA=,E25.14,3H N=,I6,3H X=,E25.14)	IBE 4240
99998 FORMAT (51H0IBESS CALLED WITH THE ORDER, ALPHA, LESS THAN ZERO	IBE 4250
* /6H KODE=,I2,7H ALPHA=,E25.14,3H N=,I6,3H X=,E25.14)	IBE 4260
99997 FORMAT (34H0IBESS CALLED WITH N LESS THAN ONE/6H KODE=,	IBE 4270
* I2,7H ALPHA=,E25.14,3H N=,I6,3H X=,E25.14)	IBE 4280
99996 FORMAT (35H0IBESS CALLED WITH X LESS THAN ZERO/6H KODE=,	IBE 4290
* I2,7H ALPHA=,E25.14,3H N=,I6,3H X=,E25.14)	IBE 4300
99995 FORMAT (42H0OVERFLOW IN IBESS, X TOO LARGE FOR KODE=1/6H KODE=,	IBE 4310
* I2,7H ALPHA=,E25.14,3H N=,I6,3H X=,E25.14)	IBE 4320
99994 FORMAT (25H0UNDERFLOW IN IBESS, LAST,I6,20H VALUE(S) OF Y ARRAY,	IBE 4330
* 17H WERE SET TO ZERO/6H KODE=,I2,7H ALPHA=,E25.14,3H N=,I6,	IBE 4340
* 3H X=,E25.14)	IBE 4350
END	IBE 4360
SUBROUTINE JAIRY(X, RX, C, AI, DAI)	JAI 10
C	JAI 20
C CDC 6600 ROUTINE	JAI 30
C 1-2-74	JAI 40
C	JAI 50
C	JAI 60
C JAIRY COMPUTES THE AIRY FUNCTION AI(X)	JAI 70
C AND ITS DERIVATIVE DAI(X) FOR JBESS	JAI 80
C INPUT	JAI 80

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C
C          X - ARGUMENT, COMPUTED BY JBESS, X.LE.(ELIM2*1.5)**(2./3.)      JAI  90
C          RX - RX=SQRT(ABS(X)), COMPUTED BY JBESS                          JAI 100
C          C - C=2.*(ABS(X)**1.5)/3., COMPUTED BY JBESS                     JAI 110
C
C          C - C=2.*(ABS(X)**1.5)/3., COMPUTED BY JBESS                     JAI 120
C
C          OUTPUT                                                            JAI 130
C
C          AI - VALUE OF FUNCTION AI(X)                                       JAI 140
C          DAI - VALUE OF THE DERIVATIVE DAI(X)                              JAI 150
C
C          AI - VALUE OF FUNCTION AI(X)                                       JAI 160
C          DAI - VALUE OF THE DERIVATIVE DAI(X)                              JAI 170
C
C          WRITTEN BY                                                         JAI 180
C
C          D. E. AMOS                                                         JAI 190
C          S. L. DANIEL                                                       JAI 200
C          M. K. WESTON                                                       JAI 210
C
C          DIMENSION AJP(19), AJN(19), A(15), B(15)                         JAI 220
C          DIMENSION AK1(14), AK2(23), AK3(14)                               JAI 230
C          DIMENSION DAJP(19), DAJN(19), DA(15), DB(15)                    JAI 240
C          DIMENSION DAK1(14), DAK2(24), DAK3(14)                           JAI 250
C          DATA N1, N2, N3, N4 /14,23,19,15/                               JAI 260
C          DATA M1, M2, M3, M4 /12,21,17,13/                               JAI 270
C          DATA FP112,CON2,CON3,CON4,CON5 / 1.30899693899575E+00,JAI 280
1 5.03154716196777E+00, 3.80004589867293E-01, 8.33333333333333E-01,JAI 290
2 8.66025403784439E-01/ JAI 300
C          DATA AK1 / 2.20423090987793E-01,-1.25290242787700E-01,JAI 310
1 1.03881163359194E-02, 8.22844152006343E-04,-2.34614345891226E-04,JAI 320
2 1.63824280172116E-05, 3.06902589573189E-07,-1.29621999359332E-07,JAI 330
3 8.22908158823668E-09, 1.53963968623298E-11,-3.39165465615682E-11,JAI 340
4 2.03253257423626E-12,-1.10679546097884E-14,-5.16169497785080E-15/JAI 350
C          DATA AK2 / 2.74366150869598E-01, 5.39790969736903E-03,JAI 360
1-1.57339220621190E-03, 4.27427528248750E-04,-1.12124917399925E-04,JAI 370
2 2.88763171318904E-05,-7.36804225370554E-06, 1.87290209741024E-06,JAI 380
3-4.75892793962291E-07, 1.21130416955909E-07,-3.09245374270614E-08,JAI 390
4 7.92454705282654E-09,-2.03902447167914E-09, 5.26863056595742E-10,JAI 400
5-1.36704767639569E-10, 3.56141039013708E-11,-9.31388296548430E-12,JAI 410
6 2.44464450473635E-12,-6.43840261990955E-13, 1.70106030559349E-13,JAI 420
7-4.50760104503281E-14, 1.19774799164811E-14,-3.19077040865066E-15/JAI 430
C          DATA AK3 / 2.80271447340791E-01,-1.78127042844379E-03,JAI 440
1 4.03422579628999E-05,-1.63249965269003E-06, 9.21181482476768E-08,JAI 450
2-6.52294330229155E-09, 5.47138404576546E-10,-5.24408251800260E-11,JAI 460
3 5.60477904117209E-12,-6.56375244639313E-13, 8.31285761966247E-14,JAI 470
4-1.12705134691063E-14, 1.62267976598129E-15,-2.46480324312426E-16/JAI 480
C          DATA AJP / 7.78952966437581E-02,-1.84356363456801E-01,JAI 490
1 3.01412605216174E-02, 3.05342724277608E-02,-4.95424702513079E-03,JAI 500
2-1.72749552563952E-03, 2.43137637839190E-04, 5.04564777517082E-05,JAI 510
3-6.16316582695208E-06,-9.03986745510768E-07, 9.70243778355884E-08,JAI 520
4 1.09639453305205E-08,-1.04716330588766E-09,-9.60359441344646E-11,JAI 530
5 8.25358789454134E-12, 6.36123439018768E-13,-4.96629614116015E-14,JAI 540
6-3.29810288929615E-15, 2.35798252031104E-16/ JAI 550
C          DATA AJN / 3.80497887617242E-02,-2.45319541845546E-01,JAI 560
1 1.65820623702696E-01, 7.49330045818789E-02,-2.63476288106641E-02,JAI 570
2-5.92535597304981E-03, 1.44744409589804E-03, 2.18311831322215E-04,JAI 580
3-4.10662077680304E-05,-4.66874994171766E-06, 7.15218807277160E-07,JAI 590
4 6.52964770854633E-08,-8.44284027565946E-09,-6.44186158976978E-10,JAI 600
5 7.20802286505285E-11, 4.72465431717846E-12,-4.66022632547045E-13,JAI 610
6-2.67762710389189E-14, 2.36161316570019E-15/ JAI 620
C          DATA A / 4.90275424742791E-01, 1.57647277946204E-03,JAI 630
1-9.66195963140306E-05, 1.35916080268815E-07, 2.98157342654859E-07,JAI 640
2-1.86824767559979E-08,-1.03685737667141E-09, 3.28660818434328E-10,JAI 650
3-2.57091410632780E-11,-2.32357655300677E-12, 9.57523279048255E-13,JAI 660
4-1.20340828049719E-13,-2.90907716770715E-15, 4.55656454580149E-15,JAI 670
5-9.99003874810259E-16/ JAI 680
C          DATA B / 2.78593552803079E-01,-3.52915691882584E-03,JAI 690
1-2.31149677384994E-05, 4.71317842263560E-06,-1.12415907931333E-07,JAI 700
2-2.00100301184339E-08, 2.60948075302193E-09,-3.55098136101216E-11,JAI 710
3-3.50849978423875E-11, 5.83007187954202E-12,-2.04644828753326E-13,JAI 720
4-1.10529179476742E-13, 2.87724778038775E-14,-2.882051110009939E-15,JAI 730
5-3.32656311696166E-16/ JAI 740
C          DATA N1D, N2D, N3D, N4D /14,24,19,15/                               JAI 750
C          DATA M1D, M2D, M3D, M4D /12,22,17,13/                               JAI 760
C          DATA DAK1 / 2.04567842307887E-01,-6.61322739905664E-02,JAI 770
1-8.49845800989287E-03, 3.12183491556289E-03,-2.70016489829432E-04,JAI 780
2-6.35636298679387E-06, 3.02397712409509E-06,-2.1831195330088E-07,JAI 790
3-5.36194289332826E-10, 1.13098035622310E-09,-7.43023834629073E-11,JAI 800
4 4.28804170826891E-13, 2.23810925754539E-13,-1.39140135641182E-14/JAI 810

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DATA DAK2 / 2.93332343883230E-01,-8.06196784743112E-03, JAI 850
1 2.42540172333140E-03,-6.82297548850235E-04, 1.85786427751181E-04, JAI 860
2-4.97457447684059E-05, 1.32090681239497E-05,-3.49528240444943E-06, JAI 870
3 9.24362451078835E-07,-2.44732671521867E-07, 6.49307837648910E-08, JAI 880
4-1.72717621501538E-08, 4.60725763604656E-09,-1.23249055291550E-09, JAI 890
5 3.30620409488102E-10,-8.89252099772401E-11, 2.39773319878298E-11, JAI 900
6-6.48013921153450E-12, 1.75510132023731E-12,-4.76303829833637E-13, JAI 910
7 1.29498241100810E-13,-3.52679622210430E-14, 9.62005151585923E-15, JAI 920
8-2.62786914342292E-15/ JAI 930
DATA DAK3 / 2.84675828811349E-01, 2.53073072619080E-03, JAI 940
1-4.83481130337976E-05, 1.84907283946343E-06,-1.01418491178576E-07, JAI 950
2 7.05925634457153E-09,-5.85325291400382E-10, 5.56357688831339E-11, JAI 960
3-5.90889094779500E-12, 6.88574353784436E-13,-8.68588256452194E-14, JAI 970
4 1.17374762617213E-14,-1.68523146510923E-15, 2.55374773097056E-16/ JAI 980
DATA DAJP / 6.53219131311457E-02,-1.20262933688823E-01, JAI 990
1 9.78010236263823E-03, 1.67948429230505E-02,-1.97146140182132E-03, JAI 1000
2-8.45560295098867E-04, 9.42889620701976E-05, 2.25827860945475E-05, JAI 1010
3-2.29067870915987E-06,-3.76343991136919E-07, 3.45663933559565E-08, JAI 1020
4 4.29611332003007E-09,-3.58673691214989E-10,-3.57245881361895E-11, JAI 1030
5 2.72696091066336E-12, 2.26120653095771E-13,-1.58763205238303E-14, JAI 1040
6-1.12604374485125E-15, 7.31327529515367E-17/ JAI 1050
DATA DAJN / 1.08594539632967E-02, 8.53313194857091E-02, JAI 1060
1-3.15277068113058E-01,-8.78420725294257E-02, 5.53251906976048E-02, JAI 1070
2 9.41674060503241E-03,-3.32187026018996E-03,-4.11157343156826E-04, JAI 1080
3 1.01297326891346E-04, 9.87633682208396E-06,-1.87312969812393E-06, JAI 1090
4-1.50798500131468E-07, 2.32687669525394E-08, 1.59599917419225E-09, JAI 1100
5-2.07665922668385E-10,-1.24103350500302E-11, 1.39631765331043E-12, JAI 1110
6 7.39400971155740E-14,-7.32887475627500E-15/ JAI 1120
DATA DA / 4.91627321104601E-01, 3.11164930427489E-03, JAI 1130
1 8.23140762854081E-05,-4.61769776172142E-06,-6.13158880534626E-08, JAI 1140
2 2.87295804656520E-08,-1.81959715372117E-09,-1.44752826642035E-10, JAI 1150
3 4.53724043420422E-11,-3.99655065847223E-12,-3.24089119830323E-13, JAI 1160
4 1.62098952568741E-13,-2.40765247974057E-14, 1.69384811284491E-16, JAI 1170
5 8.17900786477396E-16/ JAI 1180
DATA DB /-2.77571356944231E-01, 4.44212833419920E-03, JAI 1190
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2-6.24286894709776E-09,-2.36377836844577E-09, 3.16991042656673E-10, JAI 1210
3-4.40995691658191E-12,-5.18674221093575E-12, 9.64874015137022E-13, JAI 1220
4-4.90190576608710E-14,-1.77253430678112E-14, 5.55950610442662E-15, JAI 1230
5-7.11793337579530E-16/ JAI 1240
IF (X.LT.0.) GO TO 90 JAI 1250
IF (C.GT.5.) GO TO 60 JAI 1260
IF (X.GT.1.2) GO TO 30 JAI 1270
T = (X+X-1.2)*CON4 JAI 1280
TT = T + T JAI 1290
J = N1 JAI 1300
F1 = AK1(J) JAI 1310
F2 = 0. JAI 1320
DO 10 I=1,M1 JAI 1330
J = J - 1 JAI 1340
TEMP1 = F1 JAI 1350
F1 = TT*F1 - F2 + AK1(J) JAI 1360
F2 = TEMP1 JAI 1370
10 CONTINUE JAI 1380
AI = T*F1 - F2 + AK1(1) JAI 1390
J = N1D JAI 1400
F1 = DAK1(J) JAI 1410
F2 = 0. JAI 1420
DO 20 I=1,M1D JAI 1430
J = J - 1 JAI 1440
TEMP1 = F1 JAI 1450
F1 = TT*F1 - F2 + DAK1(J) JAI 1460
F2 = TEMP1 JAI 1470
20 CONTINUE JAI 1480
DAI = -(T*F1-F2+DAK1(1)) JAI 1490
RETURN JAI 1500
30 CONTINUE JAI 1510
T = (X+X-CON2)*CON3 JAI 1520
TT = T + T JAI 1530
J = N2 JAI 1540
F1 = AK2(J) JAI 1550
F2 = 0. JAI 1560
DO 40 I=1,M2 JAI 1570
J = J - 1 JAI 1580
TEMP1 = F1 JAI 1590
F1 = TT*F1 - F2 + AK2(J) JAI 1600

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F2 = TEMP1	JAI 1610
40 CONTINUE	JAI 1620
RTRX = SQRT(RX)	JAI 1630
EC = EXP(-C)	JAI 1640
AI = EC*(T*F1-F2+AK2(1))/RTRX	JAI 1650
J = N2D	JAI 1660
F1 = DAK2(J)	JAI 1670
F2 = 0.	JAI 1680
DO 50 I=1,M2D	JAI 1690
J = J - 1	JAI 1700
TEMP1 = F1	JAI 1710
F1 = TT*F1 - F2 + DAK2(J)	JAI 1720
F2 = TEMP1	JAI 1730
50 CONTINUE	JAI 1740
DAI = -EC*(T*F1-F2+DAK2(1))*RTRX	JAI 1750
RETURN	JAI 1760
60 CONTINUE	JAI 1770
T = 10./C - 1.	JAI 1780
TT = T + T	JAI 1790
J = N1	JAI 1800
F1 = AK3(J)	JAI 1810
F2 = 0.	JAI 1820
DO 70 I=1,M1	JAI 1830
J = J - 1	JAI 1840
TEMP1 = F1	JAI 1850
F1 = TT*F1 - F2 + AK3(J)	JAI 1860
F2 = TEMP1	JAI 1870
70 CONTINUE	JAI 1880
RTRX = SQRT(RX)	JAI 1890
EC = EXP(-C)	JAI 1900
AI = EC*(T*F1-F2+AK3(1))/RTRX	JAI 1910
J = N1D	JAI 1920
F1 = DAK3(J)	JAI 1930
F2 = 0.	JAI 1940
DO 80 I=1,M1D	JAI 1950
J = J - 1	JAI 1960
TEMP1 = F1	JAI 1970
F1 = TT*F1 - F2 + DAK3(J)	JAI 1980
F2 = TEMP1	JAI 1990
80 CONTINUE	JAI 2000
DAI = -RTRX*EC*(T*F1-F2+DAK3(1))	JAI 2010
RETURN	JAI 2020
90 CONTINUE	JAI 2030
IF (C.GT.5.) GO TO 120	JAI 2040
T = .4*C - 1.	JAI 2050
TT = T + T	JAI 2060
J = N3	JAI 2070
F1 = AJP(J)	JAI 2080
E1 = AJN(J)	JAI 2090
F2 = 0.	JAI 2100
E2 = 0.	JAI 2110
DO 100 I=1,M3	JAI 2120
J = J - 1	JAI 2130
TEMP1 = F1	JAI 2140
TEMP2 = E1	JAI 2150
F1 = TT*F1 - F2 + AJP(J)	JAI 2160
E1 = TT*E1 - E2 + AJN(J)	JAI 2170
F2 = TEMP1	JAI 2180
E2 = TEMP2	JAI 2190
100 CONTINUE	JAI 2200
AI = (T*E1-E2+AJN(1)) - X*(T*F1-F2+AJP(1))	JAI 2210
J = N3D	JAI 2220
F1 = DAJP(J)	JAI 2230
E1 = DAJN(J)	JAI 2240
F2 = 0.	JAI 2250
E2 = 0.	JAI 2260
DO 110 I=1,M3D	JAI 2270
J = J - 1	JAI 2280
TEMP1 = F1	JAI 2290
TEMP2 = E1	JAI 2300
F1 = TT*F1 - F2 + DAJP(J)	JAI 2310
E1 = TT*E1 - E2 + DAJN(J)	JAI 2320
F2 = TEMP1	JAI 2330
E2 = TEMP2	JAI 2340
110 CONTINUE	JAI 2350
DAI = X*X*(T*F1-F2+DAJP(1)) + (T*E1-E2+DAJN(1))	JAI 2360

RETURN	JAI 2370
120 CONTINUE	JAI 2380
T = 10./C - 1.	JAI 2390
TT = T + T	JAI 2400
J = N4	JAI 2410
F1 = A(J)	JAI 2420
E1 = B(J)	JAI 2430
F2 = 0.	JAI 2440
E2 = 0.	JAI 2450
DO 130 I=1,M4	JAI 2460
J = J - 1	JAI 2470
TEMP1 = F1	JAI 2480
TEMP2 = E1	JAI 2490
F1 = TT*F1 - F2 + A(J)	JAI 2500
E1 = TT*E1 - E2 + B(J)	JAI 2510
F2 = TEMP1	JAI 2520
E2 = TEMP2	JAI 2530
130 CONTINUE	JAI 2540
TEMP1 = T*F1 - F2 + A(1)	JAI 2550
TEMP2 = T*E1 - E2 + B(1)	JAI 2560
RTRX = SQRT(RX)	JAI 2570
CV = C - FPI12	JAI 2580
CCV = COS(CV)	JAI 2590
SCV = SIN(CV)	JAI 2600
AI = (TEMP1*CCV-TEMP2*SCV)/RTRX	JAI 2610
J = N4D	JAI 2620
F1 = DA(J)	JAI 2630
E1 = DB(J)	JAI 2640
F2 = 0.	JAI 2650
E2 = 0.	JAI 2660
DO 140 I=1,M4D	JAI 2670
J = J - 1	JAI 2680
TEMP1 = F1	JAI 2690
TEMP2 = E1	JAI 2700
F1 = TT*F1 - F2 + DA(J)	JAI 2710
E1 = TT*E1 - E2 + DB(J)	JAI 2720
F2 = TEMP1	JAI 2730
E2 = TEMP2	JAI 2740
140 CONTINUE	JAI 2750
TEMP1 = T*F1 - F2 + DA(1)	JAI 2760
TEMP2 = T*E1 - E2 + DB(1)	JAI 2770
E1 = CCV*CON5 + .5*SCV	JAI 2780
E2 = SCV*CON5 - .5*CCV	JAI 2790
DAI = (TEMP1*E1-TEMP2*E2)*RTRX	JAI 2800
RETURN	JAI 2810
END	JAI 2820
SUBROUTINE JBESS(ALPHA, N, X, Y)	JBE 10
C	JBE 20
C A CDC 6600 SUBROUTINE	JBE 30
C	JBE 40
C AUTHORS	JBE 50
C D.E. AMOS, S.L. DANIEL, AND M.K. WESTON	JBE 60
C SANDIA LABORATORIES	JBE 70
C JANUARY, 1975	JBE 80
C	JBE 90
C REFERENCES	JBE 100
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C	JBE 140
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C LARGE ORDERS. NPL MATHEMATICAL TABLES, VOL 6. HER MAJESTY-S	JBE 210
C STATIONERY OFFICE, LONDON, 1962.	JBE 220
C	JBE 230
C OLVER, F.W.J. THE ASYMPTOTIC EXPANSION OF BESSEL FUNCTIONS OF	JBE 240
C LARGE ORDER. PHIL. TRANS. A, 247, PP. 328-368, 1954.	JBE 250
C	JBE 260
C ABSTRACT	JBE 270
C JBESS COMPUTES AN N MEMBER SEQUENCE OF J BESSEL FUNCTIONS	JBE 280


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C      J/SUB(ALPHA+K-1)/(X), K=1,...,N FOR NON-NEGATIVE ALPHA AND X. JBE 290
C      A COMBINATION OF THE POWER SERIES, THE ASYMPTOTIC EXPANSION JBE 300
C      FOR X TO INFINITY AND THE UNIFORM ASYMPTOTIC EXPANSION FOR JBE 310
C      NU TO INFINITY ARE APPLIED OVER SUBDIVISIONS OF THE (NU,X) JBE 320
C      PLANE. FOR VALUES OF (NU,X) NOT COVERED BY ONE OF THESE JBE 330
C      FORMULAE, THE ORDER IS INCREMENTED OR DECREMENTED BY INTEGER JBE 340
C      VALUES INTO A REGION WHERE ONE OF THE FORMULAE APPLY. BACKWARDJBE 350
C      RECURSION IS APPLIED TO REDUCE ORDERS BY INTEGER VALUES EXCEPTJBE 360
C      WHERE THE ENTIRE SEQUENCE LIES IN THE OSCILLATORY REGION. IN JBE 370
C      THIS CASE FORWARD RECURSION IS STABLE AND VALUES FROM THE JBE 380
C      ASYMPTOTIC EXPANSION FOR X TO INFINITY START THE RECURSION JBE 390
C      WHEN IT IS EFFICIENT TO DO SO. LEADING TERMS OF THE SERIES ANDJBE 400
C      UNIFORM EXPANSION ARE TESTED FOR UNDERFLOW. IF A SEQUENCE IS JBE 410
C      REQUESTED AND THE LAST MEMBER WOULD UNDERFLOW, THE RESULT IS JBE 420
C      SET TO ZERO AND THE NEXT LOWER ORDER TRIED, ETC., UNTIL A JBE 430
C      MEMBER COMES ON SCALE OR ALL MEMBERS ARE SET TO ZERO. OVERFLOWJBE 440
C      CANNOT OCCUR. JBESS CALLS SUBROUTINE JAIRY AND FUNCTION GAMLN.JBE 450
C
C      DESCRIPTION OF ARGUMENTS JBE 470
C
C      INPUT JBE 480
C      ALPHA - ORDER OF FIRST MEMBER OF THE SEQUENCE, ALPHA.GE.0 JBE 490
C      N - NUMBER OF MEMBERS IN THE SEQUENCE, N.GE.1 JBE 510
C      X - X.GE.0 JBE 520
C
C      OUTPUT JBE 530
C      Y - A VECTOR WHOSE FIRST N COMPONENTS CONTAIN JBE 540
C      VALUES FOR J/SUB(ALPHA+K-1)/(X), K=1,...,N JBE 550
C
C      ERROR CONDITIONS JBE 560
C      IMPROPER INPUT ARGUMENTS - A FATAL ERROR JBE 570
C      UNDERFLOW - A NON-FATAL ERROR JBE 580
C
C      DOUBLE PRECISION DX, TRX, DTM, DFN JBE 590
C      DIMENSION Y(1) JBE 600
C      DIMENSION C(11,10), ALFA(26,4), BETA(26,5) JBE 610
C      DIMENSION C1(11,8), C2(11,2), ALFA1(26,2), ALFA2(26,2) JBE 620
C      DIMENSION BETA1(26,2), BETA2(26,2), BETA3(26,1) JBE 630
C      DIMENSION GAMA(26), TEMP(3), KMAX(5), AR(8), BR(10), UPOL(10) JBE 640
C      DIMENSION FNULIM(2), PP(4) JBE 650
C      DIMENSION CR(10), DR(10) JBE 660
C      EQUIVALENCE (C(1,1),C1(1,1)) JBE 670
C      EQUIVALENCE (C(1,9),C2(1,1)) JBE 680
C      EQUIVALENCE (ALFA(1,1),ALFA1(1,1)) JBE 690
C      EQUIVALENCE (ALFA(1,3),ALFA2(1,1)) JBE 700
C      EQUIVALENCE (BETA(1,1),BETA1(1,1)) JBE 710
C      EQUIVALENCE (BETA(1,3),BETA2(1,1)) JBE 720
C      EQUIVALENCE (BETA(1,5),BETA3(1,1)) JBE 730
C      DATA ELIM1,ELIM2,TOL / 667. , 644. , 1.E-15 /JBE 740
C
C      DATA PP / 8.72909153935547E+00, 2.65693932265030E-01, JBE 750
C      1 1.24578576865586E-01, 7.70133747430388E-04/ JBE 760
C TOLS=LN(1.E-3) JBE 770
C      DATA TOLS /-6.90775527898214E+00/ JBE 780
C      DATA CON1,CON2,CON548/ 6.66666666666667E-01, 3.33333333333333E-01, JBE 790
C      1 1.041666666666667E-01/ JBE 800
C      DATA RTWO,PDF,RTTP,PIDT / 1.34839972492648E+00, JBE 810
C      1 7.85398163397448E-01, 7.97884560802865E-01, 1.57079632679490E+00/ JBE 820
C      DATA FNULIM / 100. , 60. /JBE 830
C CE=-ALOG(TOL) , TCE=-0.75*ALOG(TOL) JBE 840
C      DATA CE , TCE / 3.45387763949107E+01, 2.59040822961830E+01/ JBE 850
C      DATA INLIM / 150 / JBE 860
C      DATA AR / 8.35503472222222E-02, 1.28226574556327E-01, JBE 870
C      1 2.91849026464140E-01, 8.81627267443758E-01, 3.32140828186277E+00, JBE 880
C      2 1.49957629868626E+01, 7.89230130115865E+01, 4.74451538868264E+02/ JBE 890
C      DATA BR /-1.45833333333333E-01,-9.87413194444444E-02, JBE 900
C      1-1.43312053915895E-01,-3.17227202678414E-01,-9.42429147957120E-01, JBE 910
C      2-3.51120304032635E+00,-1.57272636203680E+01,-8.22814390971859E+01, JBE 920
C      3-4.92355370523671E+02,-3.31621856854797E+03/ JBE 930
C      DATA C1 /-2.08333333333333E-01, 1.25000000000000E-01, JBE 940
C      1 9(0.) , 3.34201388888889E-01,-4.01041666666667E-01, JBE 950
C      2 7.03125000000000E-02, 8(0.) , -1.02581259645062E+00, JBE 1000
C      3 1.84646267361111E+00,-8.91210937500000E-01, 7.32421875000000E-02, JBE 1010
C      4 7(0.) , 4.66958442342625E+00,-1.12070026162230E+01, JBE 1020
C      5 8.78912353515625E+00,-2.36408691406250E+00, 1.12152099609375E-01, JBE 1030
C      6 6(0.) , -2.82120725582002E+01, 8.46362176746007E+01, JBE 1040

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7-9.18182415432400E+01, 4.25349987453885E+01, -7.36879435947963E+00, JBE 1050
 8 2.27108001708984E-01, 5(0.), 2.12570130039217E+02, JBE 1060
 9-7.65252468141182E+02, 1.05999045252800E+03, -6.99579627376133E+02, JBE 1070
 A 2.18190511744212E+02, -2.64914304869516E+01, 5.72501420974731E-01, JBE 1080
 B 4(0.), -1.91945766231841E+03, 8.06172218173731E+03, JBE 1090
 C-1.35865500064341E+04, 1.16553933368645E+04, -5.30564697861340E+03, JBE 1100
 D 1.20090291321635E+03, -1.08090919788395E+02, 1.72772750258446E+00, JBE 1110
 E 3(0.), 2.02042913309661E+04, -9.69805983886375E+04, JBE 1120
 F 1.92547001232532E+05, -2.03400177280416E+05, 1.22200464983017E+05, JBE 1130
 G-4.11926549688976E+04, 7.10951430248936E+03, -4.93915304773088E+02, JBE 1140
 H 6.07404200127348E+00, 2(0.), / JBE 1150
 DATA C2 /-2.42919187900551E+05, 1.31176361466298E+06, JBE 1160
 1-2.99801591853811E+06, 3.76327129765640E+06, -2.81356322658653E+06, JBE 1170
 2 1.26836527332162E+06, -3.31645172484564E+05, 4.52187689813627E+04, JBE 1180
 3-2.49983048181121E+03, 2.43805296995561E+01, 1(0.), JBE 1190
 4 3.28446985307204E+06, -1.97068191184322E+07, 5.09526024926646E+07, JBE 1200
 5-7.41051482115327E+07, 6.63445122747290E+07, -3.75671766607634E+07, JBE 1210
 6 1.32887671664218E+07, -2.78561812808645E+06, 3.08186404612662E+05, JBE 1220
 7-1.38860897537170E+04, 1.10017140269247E+02, / JBE 1230
 DATA ALFA1 /-4.44444444444444E-03, -9.22077922077922E-04, JBE 1240
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 2 2.6599558346255E-04, 2.61824297061501E-04, 2.48730437344656E-04, JBE 1260
 3 2.32721040083232E-04, 2.16362485712365E-04, 2.00738858762752E-04, JBE 1270
 4 1.86267636637545E-04, 1.73060775917876E-04, 1.61091705929016E-04, JBE 1280
 5 1.50274774160908E-04, 1.40503497391270E-04, 1.31668816545923E-04, JBE 1290
 6 1.23667445598253E-04, 1.16405271474738E-04, 1.09798298372713E-04, JBE 1300
 7 1.03772410422993E-04, 9.82626078369363E-05, 9.32120517249503E-05, JBE 1310
 8 8.85710852478712E-05, 8.42963105715700E-05, 8.03497548407791E-05, JBE 1320
 9 6.93735541354589E-04, 2.32241745182922E-04, -1.41986273556691E-05, JBE 1330
 A-1.16444931672049E-04, -1.50803558053049E-04, -1.55121924918096E-04, JBE 1340
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 C-1.06184319207974E-04, -9.37699549891194E-05, -8.26923045588193E-05, JBE 1360
 D-7.29374348155221E-05, -6.44042357721016E-05, -5.69611566009369E-05, JBE 1370
 E-5.04731044303562E-05, -4.48134868008883E-05, -3.98688727717599E-05, JBE 1380
 F-3.55400532972042E-05, -3.17414256609022E-05, -2.83996793904175E-05, JBE 1390
 G-2.54522720634871E-05, -2.28459297164725E-05, -2.05352753106481E-05, JBE 1400
 H-1.84816217627666E-05, -1.66519330021394E-05, / JBE 1410
 DATA ALFA2 /-3.54211971457744E-04, -1.56161263945159E-04, JBE 1420
 1 3.04465503594936E-05, 1.30198655773243E-04, 1.67471106699712E-04, JBE 1430
 2 1.70222587683593E-04, 1.56501427608595E-04, 1.36339170977445E-04, JBE 1440
 3 1.14886692029825E-04, 9.45869093034688E-05, 7.64498419250898E-05, JBE 1450
 4 6.07570334965197E-05, 4.74394299290509E-05, 3.62757512005344E-05, JBE 1460
 5 2.69939714979225E-05, 1.93210938247939E-05, 1.30056674793963E-05, JBE 1470
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 7-2.653967696797939E-06, -4.91346867098486E-06, -6.72739296091248E-06, JBE 1490
 8-8.17269379678658E-06, -9.31304715093561E-06, -1.02011418798016E-05, JBE 1500
 9 3.78194199201773E-04, 2.02471952761816E-04, -6.37938506318862E-05, JBE 1510
 A-2.38598230603006E-04, -3.10916256027362E-04, -3.13680115247576E-04, JBE 1520
 B-2.789506273791323E-04, -2.28564082619141E-04, -1.75245280340847E-04, JBE 1530
 C-1.25544063060690E-04, -8.22982872820208E-05, -4.62860730588116E-05, JBE 1540
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 E 3.62642745856794E-05, 4.58006124490189E-05, 5.24595294959114E-05, JBE 1560
 F 5.68396208545815E-05, 5.94349820393104E-05, 6.06478527578422E-05, JBE 1570
 G 6.08023907788436E-05, 6.01577894539460E-05, 5.89199657344698E-05, JBE 1580
 H 5.7251582377593E-05, 5.52804375585853E-05, / JBE 1590
 DATA BETA1 / 1.79988721413553E-02, 5.59964911064388E-03, JBE 1600
 1 2.88501402231133E-03, 1.80096606761054E-03, 1.24753110589199E-03, JBE 1610
 2 9.22878876572938E-04, 7.14430421727287E-04, 5.71787281789705E-04, JBE 1620
 3 4.69431007606482E-04, 3.93232835462917E-04, 3.34818889318298E-04, JBE 1630
 4 2.88952148495752E-04, 2.52211615549573E-04, 2.2280580798883E-04, JBE 1640
 5 1.97541838033063E-04, 1.76836855019718E-04, 1.59316899661821E-04, JBE 1650
 6 1.44347930197334E-04, 1.31448068119965E-04, 1.20245444949303E-04, JBE 1660
 7 1.10449144504599E-04, 1.01828770740567E-04, 9.41998224204238E-05, JBE 1670
 8 8.74130545753834E-05, 8.13466262162801E-05, 7.59002269646219E-05, JBE 1680
 9-1.49282953213429E-03, -8.78204709546389E-04, -5.02916549572035E-04, JBE 1690
 A-2.94822138512746E-04, -1.75463996970783E-04, -1.04008550460816E-04, JBE 1700
 B-5.96141953046458E-05, -3.12038929076098E-05, -1.26089735980230E-05, JBE 1710
 C-2.42892608575730E-07, 8.05996165414274E-06, 1.36507009262147E-05, JBE 1720
 D 1.73964125472926E-05, 1.98672978842134E-05, 2.14463263790823E-05, JBE 1730
 E 2.23954659232457E-05, 2.28967783814713E-05, 2.30785389811178E-05, JBE 1740
 F 2.30321976080909E-05, 2.28236073720349E-05, 2.25005881105292E-05, JBE 1750
 G 2.20981015361991E-05, 2.16418427448104E-05, 2.11507649256221E-05, JBE 1760
 H 2.06388749782171E-05, 2.01165241997082E-05, / JBE 1770
 DATA BETA2 / 5.52213076721293E-04, 4.47932581552385E-04, JBE 1780
 1 2.79520653992021E-04, 1.52468156198447E-04, 6.93271105657044E-05, JBE 1790
 2 1.76258683069991E-05, -1.35744996343269E-05, -3.17972413350427E-05, JBE 1800

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3-4.18861861696693E-05,-4.69004889379141E-05,-4.87665447413787E-05,JBE 1810
4-4.87010031186735E-05,-4.74755620890087E-05,-4.55813058138628E-05,JBE 1820
5-4.33309644511266E-05,-4.09230193157750E-05,-3.84822638603221E-05,JBE 1830
6-3.60857167535411E-05,-3.37793306123367E-05,-3.15888560772110E-05,JBE 1840
7-2.95269561750807E-05,-2.75978914828336E-05,-2.58006174666884E-05,JBE 1850
8-2.41308356761280E-05,-2.25823509518346E-05,-2.11479656768913E-05,JBE 1860
9-4.74617796559960E-04,-4.77864567147321E-04,-3.20390228067038E-04,JBE 1870
A-1.61105016119962E-04,-4.25778101285435E-05, 3.44571294294968E-05,JBE 1880
B 7.97092684075675E-05, 1.03138236708272E-04, 1.12466775262204E-04,JBE 1890
C 1.13103642108481E-04, 1.08651634848774E-04, 1.01437951597662E-04,JBE 1900
D 9.29298396593364E-05, 8.40293133016090E-05, 7.52727991349134E-05,JBE 1910
E 6.69632521975731E-05, 5.92564547323195E-05, 5.22169308826976E-05,JBE 1920
F 4.58539485165361E-05, 4.01445513891487E-05, 3.50481730031328E-05,JBE 1930
G 3.05157995034347E-05, 2.64956119950516E-05, 2.29363633690998E-05,JBE 1940
H 1.97893056664022E-05, 1.70091984636413E-05/ JBE 1950
  DATA BETA3 / 7.36465810572578E-04, 8.72790805146194E-04,JBE 1960
1 6.22614862573135E-04, 2.85998154194304E-04, 3.84737672879366E-06,JBE 1970
2-1.87906003636972E-04,-2.97603646594555E-04,-3.45998126832656E-04,JBE 1980
3-3.53382470916038E-04,-3.35715635775049E-04,-3.04321124789040E-04,JBE 1990
4-2.66722723047613E-04,-2.27654214122820E-04,-1.89922611854562E-04,JBE 2000
5-1.55058918599094E-04,-1.23778240761874E-04,-9.62926147717644E-05,JBE 2010
6-7.25178327714425E-05,-5.22070028895634E-05,-3.50347750511901E-05,JBE 2020
7-2.0648976103552E-05,-8.70106096849767E-06, 1.13698686675100E-06,JBE 2030
8 9.16426474122779E-06, 1.56477785428873E-05, 2.08223629482467E-05/JBE 2040
  DATA GAMA / 6.29960524947437E-01, 2.51984209978975E-01,JBE 2050
1 1.54790300415656E-01, 1.10713062416159E-01, 8.57309395527395E-02,JBE 2060
2 6.97161316958684E-02, 5.86085671893714E-02, 5.04698873536311E-02,JBE 2070
3 4.42600580689155E-02, 3.93720661543510E-02, 3.54283195924455E-02,JBE 2080
4 3.21818857502098E-02, 2.94646240791158E-02, 2.71581677112934E-02,JBE 2090
5 2.51768272973862E-02, 2.34570755306079E-02, 2.19508390134907E-02,JBE 2100
6 2.06210828235646E-02, 1.94388240897881E-02, 1.83810633800683E-02,JBE 2110
7 1.74293213231963E-02, 1.65685837786612E-02, 1.57865285987918E-02,JBE 2120
8 1.50729501494096E-02, 1.44193250839955E-02, 1.38184805735342E-02/JBE 2130
C TEST INPUT ARGUMENTS JBE 2140
  KT = 1 JBE 2150
  IF (N-1) 710, 10, 20 JBE 2160
10 KT = 2 JBE 2170
20 NN = N JBE 2180
  IF (X) 720, 30, 80 JBE 2190
30 IF (ALPHA) 700, 40, 50 JBE 2200
40 Y(1) = 1. JBE 2210
  IF (N.EQ.1) RETURN JBE 2220
  I1 = 2 JBE 2230
  GO TO 60 JBE 2240
50 I1 = 1 JBE 2250
60 DO 70 I=I1,N JBE 2260
  Y(I) = 0. JBE 2270
70 CONTINUE JBE 2280
  RETURN JBE 2290
80 CONTINUE JBE 2300
  IF (ALPHA.LT.0.) GO TO 700 JBE 2310
  DFN = DBLE(FLOAT(N)) + DBLE(ALPHA) - 1.D+0 JBE 2320
  FNU = DFN JBE 2330
  XO2 = X*.5 JBE 2340
  SXO2 = XO2*XO2 JBE 2350
C DECISION TREE FOR REGION WHERE SERIES, ASYMPTOTIC EXPANSION FOR X JBE 2360
C TO INFINITY AND ASYMPTOTIC EXPANSION FOR NU TO INFINITY ARE JBE 2370
C APPLIED. JBE 2380
  IF (SXO2.LE.(FNU+1.)) GO TO 90 JBE 2390
  TA = AMAX1(20.,FNU) JBE 2400
  IF (X.GT.TA) GO TO 120 JBE 2410
  IF (X.GT.12.) GO TO 110 JBE 2420
  XO2L = ALOG(XO2) JBE 2430
  NS = SXO2 - FNU JBE 2440
  GO TO 100 JBE 2450
90 FN = FNU JBE 2460
  FNP1 = FN + 1. JBE 2470
  XO2L = ALOG(XO2) JBE 2480
  IS = KT JBE 2490
  IF (X.LE.0.5) GO TO 330 JBE 2500
  NS = 0 JBE 2510
100 DFN = DFN + DBLE(FLOAT(NS)) JBE 2520
  FN = DFN JBE 2530
  FNP1 = FN + 1. JBE 2540
  IS = KT JBE 2550
  IF (N-1+NS.GT.0) IS = 3 JBE 2560

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GO TO 330	JBE 2570
110 NS = AMAX1(36.-FNU,0.)	JBE 2580
DFN = DFN + DBLE(FLOAT(NS))	JBE 2590
FN = DFN	JBE 2600
IS = KT	JBE 2610
IF (N-1+NS.GT.0) IS = 3	JBE 2620
GO TO 130	JBE 2630
120 CONTINUE	JBE 2640
RTX = SQRT(X)	JBE 2650
TAU = RTWO*RTX	JBE 2660
TA = TAU + FNULIM(KT)	JBE 2670
IF (FNU.LE.TA) GO TO 480	JBE 2680
FN = FNU	JBE 2690
IS = KT	JBE 2700
C UNIFORM ASYMPTOTIC EXPANSION FOR NU TO INFINITY	JBE 2710
130 CONTINUE	JBE 2720
XX = X/FN	JBE 2730
W2 = 1. - XX*XX	JBE 2740
ABW2 = ABS(W2)	JBE 2750
RA = SQRT(ABW2)	JBE 2760
IF (ABW2.GT.0.2775) GO TO 220	JBE 2770
C CASES NEAR X=FN, ABS(1.-(X/FN)**2).LE.0.2775	JBE 2780
C COEFFICIENTS OF ASYMPTOTIC EXPANSION BY SERIES	JBE 2790
C ZETA AND TRUNCATION FOR A(ZETA) AND B(ZETA) SERIES	JBE 2800
C KMAX IS TRUNCATION INDEX FOR A(ZETA) AND B(ZETA) SERIES=MAX(2,SA)	JBE 2810
SA = 0.	JBE 2820
IF (ABW2.EQ.0.) GO TO 140	JBE 2830
SA = TOLS/ALOG(ABW2)	JBE 2840
140 SB = SA	JBE 2850
DO 150 I=1,5	JBE 2860
KMAX(I) = AMAX1(SA,2.)	JBE 2870
SA = SA + SB	JBE 2880
150 CONTINUE	JBE 2890
KB = KMAX(5)	JBE 2900
KLAST = KB - 1	JBE 2910
SA = GAMA(KB)	JBE 2920
DO 160 K=1,KLAST	JBE 2930
KB = KB - 1	JBE 2940
SA = SA*W2 + GAMA(KB)	JBE 2950
160 CONTINUE	JBE 2960
Z = W2*SA	JBE 2970
AZ = ABS(Z)	JBE 2980
RTZ = SQRT(AZ)	JBE 2990
FN13 = FN**CON2	JBE 3000
RTARY = RTZ*FN13	JBE 3010
ARY = -RTARY*RTARY	JBE 3020
AZ32 = AZ*RTZ*CON1	JBE 3030
ACZ = FN*AZ32	JBE 3040
IF (Z.LE.0.) GO TO 170	JBE 3050
C TEST FOR UNDERFLOW, 1.E-280=EXP(-644.), ONE WORD LENGTH	JBE 3060
C UP FROM UNDERFLOW LIMIT OF CDC 6600	JBE 3070
IF (ACZ.GT.ELIM2) GO TO 380	JBE 3080
ARY = -ARY	JBE 3090
170 PHI = SQRT(SQRT(SA+SA+SA+SA))	JBE 3100
C B(ZETA) FOR S=0	JBE 3110
KB = KMAX(5)	JBE 3120
KLAST = KB - 1	JBE 3130
SB = BETA(KB,1)	JBE 3140
DO 180 K=1,KLAST	JBE 3150
KB = KB - 1	JBE 3160
SB = SB*W2 + BETA(KB,1)	JBE 3170
180 CONTINUE	JBE 3180
KSP1 = 1	JBE 3190
FN2 = FN*FN	JBE 3200
RFN2 = 1./FN2	JBE 3210
RDEN = 1.	JBE 3220
ASUM = 1.	JBE 3230
RELB = TOL*ABS(SB)	JBE 3240
BSUM = SB	JBE 3250
DO 200 KS=1,4	JBE 3260
KSP1 = KSP1 + 1	JBE 3270
RDEN = RDEN*RFN2	JBE 3280
C A(ZETA) AND B(ZETA) FOR S=1,2,3,4	JBE 3290
KB = KMAX(5-KS)	JBE 3300
KLAST = KB - 1	JBE 3310
SA = ALFA(KB,KS)	JBE 3320

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        SB = BETA(KB,KSP1)
        DO 190 K=1,KLAST
            KB = KB - 1
            SA = SA*W2 + ALFA(KB,KS)
            SB = SB*W2 + BETA(KB,KSP1)
190    CONTINUE
        TA = SA*RDEN
        TB = SB*RDEN
        ASUM = ASUM + TA
        BSUM = BSUM + TB
        IF (ABS(TA).LE.TOL .AND. ABS(TB).LE.RELB) GO TO 210
200    CONTINUE
210    CONTINUE
        BSUM = BSUM/(FN*FN13)
        GO TO 300
220    CONTINUE
        TAU = 1./RA
        T2 = 1./W2
        IF (W2.GE.0.) GO TO 230
C CASES FOR (X/FN).GT.SQRT(1.2775)
        AZ32 = ABS(RA-ATAN(RA))
        ACZ = AZ32*FN
        CZ = -ACZ
        Z32 = 1.5*AZ32
        RTZ = Z32**CON2
        FN13 = FN**CON2
        RTARY = RTZ*FN13
        ARY = -RTARY*RTARY
        GO TO 240
230    CONTINUE
C CASES FOR (X/FN).LT.SQRT(0.7225)
        AZ32 = ABS(ALOG((1.+RA)/XX)-RA)
C TEST FOR UNDERFLOW, 1.E-280 = EXP(-644.), ONE WORD LENGTH
C UP FROM UNDERFLOW LIMIT OF CDC 6600
        ACZ = AZ32*FN
        CZ = ACZ
        IF (ACZ.GT.ELIM2) GO TO 380
        Z32 = 1.5*AZ32
        RTZ = Z32**CON2
        FN13 = FN**CON2
        RTARY = RTZ*FN13
        ARY = RTARY*RTARY
240    CONTINUE
        PHI = SQRT((RTZ+RTZ)*TAU)
        TB = 1.
        ASUM = 1.
        TFN = TAU/FN
        UPOL(1) = 1.
        UPOL(2) = (C(1,1)*T2+C(2,1))*TFN
        RCZ = CON1/CZ
        CRZ32 = CON548*RCZ
        BSUM = UPOL(2) + CRZ32
        RELB = TOL*ABS(BSUM)
        AP = TFN
        KS = 0
        KP1 = 2
        RZDEN = RCZ
        DO 280 LR=2,8,2
C COMPUTE TWO U POLYNOMIALS FOR NEXT A(ZETA) AND B(ZETA)
        LRP1 = LR + 1
        DO 260 K=LR,LRP1
            KS = KS + 1
            KP1 = KP1 + 1
            S1 = C(1,K)
            DO 250 J=2,KP1
                S1 = S1*T2 + C(J,K)
250    CONTINUE
            AP = AP*TFN
            UPOL(KP1) = AP*S1
            CR(KS) = BR(KS)*RZDEN
            RZDEN = RZDEN*RCZ
            DR(KS) = AR(KS)*RZDEN
260    CONTINUE
            SUMA = UPOL(LRP1)
            SUMB = UPOL(LR+2) + UPOL(LRP1)*CRZ32
            JU = LRP1

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JBE 3330
 JBE 3340
 JBE 3350
 JBE 3360
 JBE 3370
 JBE 3380
 JBE 3390
 JBE 3400
 JBE 3410
 JBE 3420
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 JBE 3990
 JBE 4000
 JBE 4010
 JBE 4020
 JBE 4030
 JBE 4040
 JBE 4050
 JBE 4060
 JBE 4070
 JBE 4080

DO 270 JR=1,LR	JBE 4090
JU = JU - 1	JBE 4100
SUMA = SUMA + CR(JR)*UPOL(JU)	JBE 4110
SUMB = SUMB + DR(JR)*UPOL(JU)	JBE 4120
270 CONTINUE	JBE 4130
TB = -TB	JBE 4140
IF (W2.GT.0.) TB = ABS(TB)	JBE 4150
ASUM = ASUM + SUMA*TB	JBE 4160
BSUM = BSUM + SUMB*TB	JBE 4170
IF (ABS(SUMA).LE.TOL .AND. ABS(SUMB).LE.RELB) GO TO 290	JBE 4180
280 CONTINUE	JBE 4190
290 TB = RTARY	JBE 4200
IF (W2.GT.0.) TB = -TB	JBE 4210
BSUM = BSUM/TB	JBE 4220
300 CONTINUE	JBE 4230
CALL JAIKY(ARY, RTARY, ACZ, AI, DAI)	JBE 4240
TEMP(IS) = PHI*(AI*ASUM+DAI*BSUM)/FNI3	JBE 4250
GO TO (320, 450, 610), IS	JBE 4260
310 TEMP(1) = TEMP(3)	JBE 4270
KT = 1	JBE 4280
320 IS = 2	JBE 4290
DFN = DFN - 1.D+0	JBE 4300
FN = DFN	JBE 4310
GO TO 130	JBE 4320
C SERIES FOR (X/2)**2.LE.NU+1	JBE 4330
330 CONTINUE	JBE 4340
GLN = GMLN(FNP1)	JBE 4350
ARG = FN*XO2L - GLN	JBE 4360
IF (ARG.LT.-ELIM1) GO TO 400	JBE 4370
EARG = EXP(ARG)	JBE 4380
340 CONTINUE	JBE 4390
S = 1.	JBE 4400
AK = 3.	JBE 4410
T2 = 1.	JBE 4420
T = 1.	JBE 4430
S1 = FN	JBE 4440
DO 350 K=1,17	JBE 4450
S2 = T2 + S1	JBE 4460
T = -T*SXO2/S2	JBE 4470
S = S + T	JBE 4480
IF (ABS(T).LT.TOL) GO TO 360	JBE 4490
T2 = T2 + AK	JBE 4500
AK = AK + 2.	JBE 4510
S1 = S1 + FN	JBE 4520
350 CONTINUE	JBE 4530
360 CONTINUE	JBE 4540
TEMP(IS) = S*EARG	JBE 4550
GO TO (370, 450, 600), IS	JBE 4560
370 EARG = EARG*FN/XO2	JBE 4570
DFN = DFN - 1.D+0	JBE 4580
FN = DFN	JBE 4590
IS = 2	JBE 4600
GO TO 340	JBE 4610
C SET UNDERFLOW VALUE AND UPDATE PARAMETERS	JBE 4620
380 Y(NN) = 0.	JBE 4630
NN = NN - 1	JBE 4640
DFN = DFN - 1.D+0	JBE 4650
FN = DFN	JBE 4660
IF (NN-1) 440, 390, 130	JBE 4670
390 KT = 2	JBE 4680
IS = 2	JBE 4690
GO TO 130	JBE 4700
400 Y(NN) = 0.	JBE 4710
NN = NN - 1	JBE 4720
FNP1 = FN	JBE 4730
DFN = DFN - 1.D+0	JBE 4740
FN = DFN	JBE 4750
IF (NN-1) 440, 410, 420	JBE 4760
410 KT = 2	JBE 4770
IS = 2	JBE 4780
420 IF (SXO2.LE.FNP1) GO TO 430	JBE 4790
GO TO 130	JBE 4800
430 ARG = ARG - XO2L + ALOG(FNP1)	JBE 4810
IF (ARG.LT.-ELIM1) GO TO 400	JBE 4820
GO TO 330	JBE 4830
440 NZ = N - NN	JBE 4840

PRINT 99996, NZ, ALPHA, N, X	JBE 4850
RETURN	JBE 4860
C BACKWARD RECURSION SECTION	JBE 4870
450 CONTINUE	JBE 4880
NZ = N - NN	JBE 4890
IF (NZ.NE.0) PRINT 99996, NZ, ALPHA, N, X	JBE 4900
IF (KT.EQ.2) GO TO 470	JBE 4910
C BACKWARD RECUR FROM INDEX ALPHA+NN-1 TO ALPHA	JBE 4920
Y(NN) = TEMP(1)	JBE 4930
Y(NN-1) = TEMP(2)	JBE 4940
IF (NN.EQ.2) RETURN	JBE 4950
DX = X	JBE 4960
TRX = 2.D+0/DX	JBE 4970
DTM = DFN*TRX	JBE 4980
TM = DTM	JBE 4990
K = NN + 1	JBE 5000
DO 460 I=3,NN	JBE 5010
K = K - 1	JBE 5020
Y(K-2) = TM*Y(K-1) - Y(K)	JBE 5030
DTM = DTM - TRX	JBE 5040
TM = DTM	JBE 5050
460 CONTINUE	JBE 5060
RETURN	JBE 5070
470 Y(1) = TEMP(2)	JBE 5080
RETURN	JBE 5090
C ASYMPTOTIC EXPANSION FOR X TO INFINITY WITH FORWARD RECURSION IN	JBE 5100
C OSCILLATORY REGION X.GT.MAX(20, NU), PROVIDED THE LAST MEMBER	JBE 5110
C OF THE SEQUENCE IS ALSO IN THE REGION.	JBE 5120
480 CONTINUE	JBE 5130
IN = ALPHA - TAU + 2.	JBE 5140
IF (IN.LE.0) GO TO 490	JBE 5150
INP1 = IN + 1	JBE 5160
DALPHA = ALPHA - FLOAT(INP1)	JBE 5170
KT = 1	JBE 5180
GO TO 500	JBE 5190
490 DALPHA = ALPHA	JBE 5200
IN = 0	JBE 5210
500 IS = KT	JBE 5220
ARG = X - PIDT*DALPHA - PDF	JBE 5230
SA = SIN(ARG)	JBE 5240
SB = COS(ARG)	JBE 5250
RA = RTTP/RTX	JBE 5260
ETX = 8.*X	JBE 5270
510 DX = DALPHA	JBE 5280
DX = DX + DX	JBE 5290
DTM = DX*DX	JBE 5300
T2 = DTM - 1.D+0	JBE 5310
T2 = T2/ETX	JBE 5320
S2 = T2	JBE 5330
RELB = TOL*ABS(T2)	JBE 5340
T1 = ETX	JBE 5350
S1 = 1.	JBE 5360
FN = 1.	JBE 5370
AK = 8.	JBE 5380
DO 520 K=1,13	JBE 5390
T1 = T1 + ETX	JBE 5400
FN = FN + AK	JBE 5410
DX = FN	JBE 5420
TRX = DTM - DX	JBE 5430
AP = TRX	JBE 5440
T2 = -T2*AP/T1	JBE 5450
S1 = S1 + T2	JBE 5460
T1 = T1 + ETX	JBE 5470
AK = AK + 8.	JBE 5480
FN = FN + AK	JBE 5490
DX = FN	JBE 5500
TRX = DTM - DX	JBE 5510
AP = TRX	JBE 5520
T2 = T2*AP/T1	JBE 5530
S2 = S2 + T2	JBE 5540
IF (ABS(T2).LE.RELB) GO TO 530	JBE 5550
AK = AK + 8.	JBE 5560
520 CONTINUE	JBE 5570
530 TEMP(IS) = RA*(S1*SB-S2*SA)	JBE 5580
GO TO (540, 550), IS	JBE 5590
540 DALPHA = DALPHA + 1.	JBE 5600

IS = 2	JBE 5610
TB = SA	JBE 5620
SA = -SB	JBE 5630
SB = TB	JBE 5640
GO TO 510	JBE 5650
C FORWARD RECURSION SECTION	JBE 5660
550 IF (KT.EQ.2) GO TO 470	JBE 5670
S1 = TEMP(1)	JBE 5680
S2 = TEMP(2)	JBE 5690
TX = 2./X	JBE 5700
TM = DALPHA*TX	JBE 5710
IF (IN.EQ.0) GO TO 570	JBE 5720
C FORWARD RECUR TO INDEX ALPHA	JBE 5730
DO 560 I=1,IN	JBE 5740
S = S2	JBE 5750
S2 = TM*S2 - S1	JBE 5760
TM = TM + TX	JBE 5770
S1 = S	JBE 5780
560 CONTINUE	JBE 5790
IF (NN.EQ.1) GO TO 590	JBE 5800
S = S2	JBE 5810
S2 = TM*S2 - S1	JBE 5820
TM = TM + TX	JBE 5830
S1 = S	JBE 5840
570 CONTINUE	JBE 5850
C FORWARD RECUR FROM INDEX ALPHA TO ALPHA+N-1	JBE 5860
Y(1) = S1	JBE 5870
Y(2) = S2	JBE 5880
IF (NN.EQ.2) RETURN	JBE 5890
DO 580 I=3,NN	JBE 5900
Y(I) = TM*Y(I-1) - Y(I-2)	JBE 5910
TM = TM + TX	JBE 5920
580 CONTINUE	JBE 5930
RETURN	JBE 5940
590 Y(1) = S2	JBE 5950
RETURN	JBE 5960
C BACKWARD RECURSION WITH NORMALIZATION BY	JBE 5970
C ASYMPTOTIC EXPANSION FOR NU TO INFINITY OR POWER SERIES.	JBE 5980
600 CONTINUE	JBE 5990
C COMPUTATION OF LAST ORDER FOR SERIES NORMALIZATION	JBE 6000
KM = AMAX1(3.-FN,0.)	JBE 6010
TFN = FN + FLOAT(KM)	JBE 6020
TA = (GLN+TFN-0.9189385332-0.0833333333/TFN)/(TFN+0.5)	JBE 6030
TA = XO2L - TA	JBE 6040
TB = -(1.-1.5/TFN)/TFN	JBE 6050
IN = CE/(-TA+SQRT(TA*TA-CE*TB)) + 1.5	JBE 6060
IN = IN + KM	JBE 6070
GO TO 650	JBE 6080
610 CONTINUE	JBE 6090
C COMPUTATION OF LAST ORDER FOR ASYMPTOTIC EXPANSION NORMALIZATION	JBE 6100
GLN = AZ32 + RA	JBE 6110
IF (ARY.GT.30.) GO TO 630	JBE 6120
RDEN = (PP(4)*ARY+PP(3))*ARY + 1.	JBE 6130
RZDEN = PP(1) + PP(2)*ARY	JBE 6140
TA = RZDEN/RDEN	JBE 6150
IF (W2.LT.0.10) GO TO 620	JBE 6160
TB = GLN/RTARY	JBE 6170
GO TO 640	JBE 6180
620 TB = (1.259921049+0.1679894730*W2)/FN13	JBE 6190
GO TO 640	JBE 6200
630 CONTINUE	JBE 6210
TA = CON1*TCE/ACZ	JBE 6220
TA = ((0.0493827160*TA-0.1111111111)*TA+0.6666666667)*TA*ARY	JBE 6230
IF (W2.LT.0.10) GO TO 620	JBE 6240
TB = GLN/RTARY	JBE 6250
640 IN = TA/TB + 1.5	JBE 6260
IF (IN.GT.INLIM) GO TO 310	JBE 6270
650 DX = FLOAT(IN)	JBE 6280
DTM = DFN + DX	JBE 6290
DX = X	JBE 6300
TRX = 2.D+0/DX	JBE 6310
DTM = DTM*TRX	JBE 6320
TM = DTM	JBE 6330
TA = 0.	JBE 6340
TB = TOL	JBE 6350
KK = 1	JBE 6360

660	CONTINUE	JBE 6370
C	BACKWARD RECUR UNINDEXED	JBE 6380
	DO 670 I=1,IN	JBE 6390
	TB = TM*TB - TA	JBE 6400
	DTM = DTM - TRX	JBE 6410
	TM = DTM	JBE 6420
670	CONTINUE	JBE 6430
C	NORMALIZATION	JBE 6440
	IF (KK.NE.1) GO TO 680	JBE 6450
	TA = (TA/TB)*TEMP(3)	JBE 6460
	TB = TEMP(3)	JBE 6470
	KK = 2	JBE 6480
	IN = NS	JBE 6490
	IF (NS.NE.0) GO TO 660	JBE 6500
680	Y(NN) = TB	JBE 6510
	NZ = N - NN	JBE 6520
	IF (NZ.NE.0) PRINT 99996, NZ, ALPHA, N, X	JBE 6530
	IF (NN.EQ.1) RETURN	JBE 6540
	TB = TM*TB - TA	JBE 6550
	DTM = DTM - TRX	JBE 6560
	TM = DTM	JBE 6570
	K = NN - 1	JBE 6580
	Y(K) = TB	JBE 6590
	IF (NN.EQ.2) RETURN	JBE 6600
	KM = K - 1	JBE 6610
C	BACKWARD RECUR INDEXED	JBE 6620
	DO 690 I=1,KM	JBE 6630
	Y(K-1) = TM*Y(K) - Y(K+1)	JBE 6640
	DTM = DTM - TRX	JBE 6650
	TM = DTM	JBE 6660
	K = K - 1	JBE 6670
690	CONTINUE	JBE 6680
	RETURN	JBE 6690
700	PRINT 99999, ALPHA, N, X	JBE 6700
	STOP	JBE 6710
710	PRINT 99998, ALPHA, N, X	JBE 6720
	STOP	JBE 6730
720	PRINT 99997, ALPHA, N, X	JBE 6740
	STOP	JBE 6750
99999	FORMAT (51H0JBESS CALLED WITH THE ORDER, ALPHA, LESS THAN ZERO	JBE 6760
	* /7H ALPHA=,E25.14,3H N=,I6,3H X=,E25.14)	JBE 6770
99998	FORMAT (34H0JBESS CALLED WITH N LESS THAN ONE/7H ALPHA=,	JBE 6780
	* E25.14,3H N=,I6,3H X=,E25.14)	JBE 6790
99997	FORMAT (35H0JBESS CALLED WITH X LESS THAN ZERO/7H ALPHA=,	JBE 6800
	* E25.14,3H N=,I6,3H X=,E25.14)	JBE 6810
99996	FORMAT (25H0UNDERFLOW IN JBESS, LAST,I6,20H VALUE(S) OF Y ARRAY,	JBE 6820
	* 17H WERE SET TO ZERO/7H ALPHA=,E25.14,3H N=,I6,3H X=,E25.14)	JBE 6830
	END	JBE 6840

ERRATUM: ALGORITHM 511

CDC 6600 Subroutines IBESS and JBESS for Bessel Functions $I_\nu(x)$ and $J_\nu(x)$, $x \geq 0$, $\nu \geq 0$ [S18]

[D.E. Amos, S.L. Daniel, and M.K. Weston, *ACM Trans. Math. Software* 3, 1 (March 1977), 93-95]

Donald E. Amos [Recd 28 July 1977]

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The following changes should be made in this algorithm:

- (1) Add (immediately after line JBE 6390)¹

S = TB JBE 6391

¹ To make this change, see "Collected Algorithms from ACM," pages 511-P1 through 511-P21.

(2) Add (immediately after line JBE 6400)¹

TA = S JBE 6401

(3) Replace the date

1977

by the date

March, 1977

in the lines IBE 180 and JBE 180,² and the date

1976

by the date

March, 1977

in the line GAM 180.¹

² To make these changes, see the issue of *ACM Trans. Math. Software* referenced above, pages 94 and 95.

ALGORITHM 512

A Normalized Algorithm for the Solution of Positive Definite Symmetric Quindagonal Systems of Linear Equations [F4]

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Key Words and Phrases: linear equations, normalized solution, periodic quindagonal, symmetric positive definite

CR Categories: 5.14, 5.17

Language: Fortran

DESCRIPTION

In the finite difference solution of elliptic partial differential equations using block iterative methods, there occurs within the iteration scheme a need to solve many subsets of the resulting set of linear equations for each block and for each iteration. The form of the coefficient matrix B of such subsets of equations depends on the particular differential equation and the region of integration and on the way in which the grid points in the finite difference scheme are ordered. In this algorithm we are concerned with the case when B is symmetric, positive definite, and of the form

$$\begin{pmatrix} c_1 & b_1 & a_1 & & & & & & & & a_{N-1} & b_N \\ b_1 & c_2 & b_2 & a_2 & & & & & & & & a_N \\ a_1 & b_2 & c_3 & b_3 & a_3 & & & & & & & \\ & a_2 & & & & & & & & & & \\ & & \cdot & \cdot & \cdot & \cdot & \cdot & & & & & \\ & & & \cdot & \cdot & \cdot & \cdot & \cdot & & & & \\ & & & & \cdot & \cdot & \cdot & \cdot & \cdot & & & \\ & & & & & a_{N-4} & b_{N-3} & c_{N-2} & b_{N-2} & a_{N-2} & & \\ a_{N-1} & & & & & a_{N-3} & b_{N-2} & c_{N-1} & b_{N-1} & & & \\ b_N & a_N & & & & a_{N-2} & b_{N-1} & c_N & & & & \end{pmatrix} \quad (1)$$

Such systems arise, for example, from a two-line grouping of the mesh points for a self-adjoint partial differential equation in a rectangular region with periodic boundary conditions in one or both coordinate directions. Similar systems also arise in the solution of the biharmonic equation in a rectangular region subject to periodicity conditions [4]. For nonsymmetric matrices such as might occur if the differential equation were not self-adjoint or for circular regions [1], an algorithm has been given [2]. For the symmetric, positive definite case, however, it is worth-

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while to consider a normalized form of the algorithm following Cuthill and Varga [3]. This approach makes use of the result:

If A is a real $N \times N$, symmetric positive definite matrix, then there exists a unique, positive, diagonal matrix D and a unique, real, upper triangular matrix T , with unit diagonal elements such that

$$A = DT'TD \tag{2}$$

where T' denotes the transpose of T .

Assume, then, that in the system of equations

$$B\phi = s \tag{3}$$

the matrix B is symmetric, positive definite, and of the form given by (1). Then B has the unique factorization

$$B = DT'TD$$

where

$$D = \text{diag} \{d_1, d_2, d_3, \dots, d_N\} \tag{4}$$

and T is given by

$$T = \begin{pmatrix} 1 & e_1 & f_1 & & & & g_1 & & h_1 \\ & 1 & e_2 & f_2 & & & g_2 & & h_2 \\ & & \cdot & \cdot & \cdot & & \cdot & & \cdot \\ & & & \cdot & \cdot & \cdot & \cdot & & \cdot \\ & & & & \cdot & \cdot & \cdot & & \cdot \\ & & & & & 1 & e_{N-4} & f_{N-4} & g_{N-4} & h_{N-4} \\ & & & & & & 1 & e_{N-3} & (f_{N-3} + g_{N-3}) & h_{N-3} \\ & & & & & & & 1 & (e_{N-2} + g_{N-2}) & (f_{N-2} + h_{N-2}) \\ & & & & & & & & 1 & (e_{N-1} + h_{N-1}) \\ & & & & & & & & & 1 \end{pmatrix}. \tag{5}$$

The elements of the matrices D and T are given by the following relations.

For the first $(N - 2)$ columns, the factorization is straightforward, as there is no interaction between the elements in T ; so that

$$d_1 = (c_1)^{1/2}, \quad d_2 = (c_2 - b_1^2/d_1^2)^{1/2}, \quad e_1 = b_1/d_1 d_2,$$

and for $i = 3, 4, \dots, N - 2$

$$u = a_{i-2}/d_{i-2}; \quad v = b_{i-1}/d_{i-1} - e_{i-2}u$$

and

$$d_i = (c_i - u^2 - v^2)^{1/2}, \quad f_{i-2} = u/d_i, \quad e_{i-1} = v/d_i.$$

For the last two columns auxiliary vectors x and y are used; so that for column $(N - 1)$ if

$$\begin{aligned} u &= a_{N-3}/d_{N-3}, & v &= b_{N-2}/d_{N-2} - e_{N-3}u, \\ x_1 &= a_{N-1}/d_1, & x_2 &= -e_1x_1, \end{aligned}$$

and for $i = 3, 4, \dots, N - 2$

$$x_i = -e_{i-1}x_{i-1} - f_{i-2}x_{i-2},$$

then

$$\begin{aligned} d_{N-1} &= \left\{ c_{N-1} - \left[\sum_{i=1}^{N-4} x_i^2 + (x_{N-3} + u)^2 + (x_{N-2} + v)^2 \right] \right\}^{1/2}, \\ f_{N-3} &= u/d_{N-1}, \end{aligned}$$

$$e_{N-2} = v/d_{N-1},$$

and for $i = 1, 2, \dots, N - 2$

$$g_i = x_i/d_{N-1}.$$

For the final column, we set

$$u = a_{N-2}/d_{N-2}, \quad v = b_{N-1}/d_{N-1} - e_{N-2}u,$$

$$y_1 = b_N/d_1, \quad y_2 = a_N/d_2 - e_1y_1,$$

and for $i = 3, 4, \dots, N - 1$

$$y_i = -e_{i-1}y_{i-1} - f_{i-2}y_{i-2},$$

and finally

$$y_{N-1} = y_{N-1} - \sum_{i=1}^{N-2} g_i y_i - g_{N-2}u.$$

Then

$$d_N = \{c_N - \sum_{i=1}^{N-3} y_i^2 - (y_{N-2} + u)^2 - (y_{N-1} + v)^2\}^{1/2},$$

$$f_{N-2} = u/d_N,$$

$$e_{N-1} = v/d_N,$$

and for $i = 1, 2, \dots, N - 1$

$$h_i = y_i/d_N.$$

The factorization stage is then complete. To solve the linear system (3), with coefficient matrix B as given by (1), we rewrite it in the equivalent form:

$$DT'TD\phi = \mathbf{s} \quad \text{i.e.} \quad T'T(D\phi) = D^{-1}\mathbf{s}; \quad (6)$$

so that putting

$$D\phi = \mathbf{z}, \quad D^{-1}\mathbf{s} = \mathbf{q}, \quad (7)$$

the system becomes

$$T'T\mathbf{z} = \mathbf{q},$$

which may be solved directly for \mathbf{z} in terms of the auxiliary vector \mathbf{p} whose elements are given by

$$p_1 = q_1, \quad p_2 = q_2 - e_1p_1,$$

and for $i = 3, 4, \dots, N - 2$

$$p_i = q_i - e_{i-1}p_{i-1} - f_{i-2}p_{i-2}.$$

Finally,

$$p_{N-1} = q_{N-1} - (e_{N-2} + g_{N-2})p_{N-2} - (f_{N-3} + g_{N-3})p_{N-3} - \sum_{i=1}^{N-4} g_i p_i,$$

$$p_N = q_N - (e_{N-1} + h_{N-1})p_{N-1} - (f_{N-2} + h_{N-2})p_{N-2} - \sum_{i=1}^{N-3} h_i p_i.$$

The elements of \mathbf{z} are then obtained by a back substitution process so that

$$z_N = p_N,$$

$$z_{N-1} = p_{N-1} - (e_{N-1} + h_{N-1})z_N,$$

$$z_{N-2} = p_{N-2} - (e_{N-2} + g_{N-2})z_{N-1} - (f_{N-2} + h_{N-2})z_N,$$

and for $i = N - 3, \dots, 3, 2, 1$

$$z_i = p_i - e_i z_{i+1} - f_i z_{i+2} - g_i z_{N-1} - h_i z_N.$$

The final solution ϕ is then obtained from $\phi = D^{-1}\mathbf{z}$ using only one division per component.

Storage Management Used in the Program

The algorithm has been written in the form of three subroutines, as follows.

1. FACTOR(A,B,C,G,H,N)

This subroutine performs the factorization stage. A matrix of the form (1) is input as three vectors \mathbf{a} , \mathbf{b} , \mathbf{c} , which are overwritten during the course of the factorization. The diagonal elements of D , i.e. d_1 to d_n , replace the coefficients c_1 to c_n ; in the vectors \mathbf{b} and \mathbf{a} are stored the elements e_1 to e_{n-1} and f_1 to f_{n-2} , respectively, which occur in the matrix T . In the factorization formulas two vectors \mathbf{x} and \mathbf{y} are used, but in the program these are replaced by \mathbf{g} and \mathbf{h} which are in turn overwritten by the g_1 to g_{n-2} and h_1 to h_{n-1} values which occur in the last two columns of T .

2. RHS(D,S,N)

This subroutine forms $D^{-1}\mathbf{s}$, where D is a diagonal matrix whose elements are stored in a vector \mathbf{d} . The input vector \mathbf{s} is overwritten by the result vector.

3. SOLVE(E,F,G,H,Q,N)

This subroutine solves the set of equations $T'T\mathbf{z} = \mathbf{q}$ for T as given by (5). The solution is effected in two stages, viz. a forward substitution for \mathbf{p} where $T'\mathbf{p} = \mathbf{q}$ and a backward substitution for \mathbf{z} where $T\mathbf{z} = \mathbf{p}$. The input vector \mathbf{q} is overwritten first by the intermediate solution \mathbf{p} and finally by the result \mathbf{z} .

To solve (3) directly with B as given by (1) would then require the subroutine calls:

```
CALL FACTOR(A,B,C,G,H,N)
CALL RHS(C,S,N)
CALL SOLVE(B,A,G,H,S,N)
CALL RHS(C,S,N)
```

The reason for the three subroutines (instead of a single one), however, is that in a typical block successive overrelaxation process the three stages need to be performed at different points and possibly a different number of times. For instance, if the matrix B remains the same for each block, then the factorization need be performed only once. Moreover, the set of equations (3) is typically solved as part of an iterative procedure so that for block r , if k represents the iteration count then we have for successive block overrelaxation

$$B_r \phi_r^{(k+1)} = \mathbf{s}(\phi_{r-1}^{(k+1)}, \phi_r^{(k)}, \phi_{r+1}^{(k)}).$$

However, if $D_r \phi_r^{(k)} = \mathbf{z}_r^{(k)}$, then the iteration is carried out in the transformed variable \mathbf{z} , so that

$$T_r' T_r \mathbf{z}_r^{(k+1)} = \mathbf{q}(\mathbf{z}_{r-1}^{(k+1)}, \mathbf{z}_r^{(k)}, \mathbf{z}_{r+1}^{(k)})$$

until for some value K of k , $\mathbf{z}_r^{(K+1)}$ and $\mathbf{z}_r^{(K)}$ agree to some specified accuracy for all values of r . The final result is then obtained from $\phi = D^{-1}\mathbf{z}$. The process is considered to be a normalized Choleski technique used in an inner loop where the need to include the division is eliminated until the final result is achieved.

Normalized Algorithm Versus Choleski Factorization

Let us briefly consider the differences between using the normalized factorization $B = DT'TD$ as opposed to a standard Choleski factorization $B = LL'$ which takes into account the structure of the nonzero elements of B .

The normalized form is clearly similar to a Choleski factorization since DT' is a lower triangular matrix, say \tilde{L} , so that $B = \tilde{L}\tilde{L}'$.

Moreover, if we consider $B = LL'$, then the formulas for determining the elements of L are very similar to those already given for the elements of D and T .

However, a work count of the two algorithms shows that the normalized version requires $2n$ more divisions than the Choleski form, this extra work occurring in the factorization stage. Having factorized B , the two solution processes for determining ϕ from $B\phi = s$ require the same amount of work. Consequently, for the straightforward solution of a set of equations $B\phi = s$ with B as in (1), the Choleski factorization is preferable because less work is involved.

However, as we have seen, the B matrices are likely to occur in problems where the solution of $B\phi = s$ is incorporated into an iterative procedure. In this case for the normalized version, n divisions per block would be saved in each iteration in excess of one (since the iteration is effected in a transformed variable thus saving the final division). Let us consider all possible cases for the B matrices, i.e.

- (i) $B = B_r^{(k)}$: the matrices change from block to block and with each iteration;
- (ii) $B = B^{(k)}$: the matrices change with each iteration but not from block to block;
- (iii) $B = B_r$: the matrices change from block to block but not with each iteration;
- (iv) $B = B$: the matrices remain constant throughout the process.

Cases (iii) and (iv) are the most likely to occur in practice.

For (i), the Choleski factorization always involves less work.

For (ii), the Choleski form is more efficient if the number of blocks is less than three. For three or more blocks, the normalized form is more efficient if the number of iterations exceeds two.

For (iii) and (iv), i.e. the most likely cases, the normalized form is more efficient, from the point of view of number of operations involved, if the number of iterations is again more than two.

Consequently we conclude that for the majority of problems involving a block overrelaxation technique in which the blocks have a coefficient matrix of the form (1), the normalized algorithm as presented here would involve less work than a sparse Choleski form.

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ALGORITHM

SUBROUTINE FACTOR(A, B, C, G, H, N)	FAC	10
C THE SUBROUTINE IS A NORMALISED FACTORISATION OF A SQUARE MATRIX OF	FAC	20
C ORDER GREATER THAN 4.	FAC	30
C THE COEFFICIENT MATRIX P IS SYMMETRIC, POSITIVE DEFINITE AND	FAC	40
C QUINDIAGONAL WITH NON-ZERO ELEMENTS ALSO IN THE LAST TWO COLUMNS OF	FAC	50
C THE FIRST ROW, THE LAST COLUMN OF THE SECOND ROW, THE FIRST COLUMN	FAC	60
C OF ROW (N-1) AND THE FIRST TWO COLUMNS OF THE LAST ROW. A, B, C ARE	FAC	70
C VECTORS EACH OF N ELEMENTS CONTAINING RESPECTIVELY THE SUPER-SUPER-	FAC	80
C DIAGONAL, SUPER-DIAGONAL AND DIAGONAL ELEMENTS OF P I.E. P(I, I+2),	FAC	90
C P(I, I+1) AND P(I, I). THE MATRIX ELEMENTS P(1, N-1), P(1, N), P(2, N)	FAC	100
C ARE STORED IN A(N-1), B(N), A(N) RESPECTIVELY, AS ARE P(N-1, 1),	FAC	110
C P(N, 1) AND P(N, 2)	FAC	120
C THE MATRIX P IS FACTORISED INTO P=DRTD WHERE D IS A DIAGONAL MATRIX	FAC	130
C AND T IS A REAL, UPPER TRIANGULAR MATRIX WITH UNIT DIAGONAL ELEMENTS	FAC	140
C AND NON-ZERO ELEMENTS IN THE LAST TWO COLUMNS AND ON THE SUPER- AND	FAC	150
C SUPER-SUPER-DIAGONALS. R DENOTES THE TRANSPOSE OF T.	FAC	160
C THE INPUT VECTORS ARE OVERRITTEN DURING THE FACTORISATION STAGE SO	FAC	170
C THAT VECTOR C CONTAINS THE DIAGONAL ELEMENTS OF D, VECTOR B CONTAINS	FAC	180

```

C THE ELEMENTS T(I,I+1) FOR I=1 TO N-3 AND VECTOR A CONTAINS THE FAC 190
C ELEMENTS T(I,I+2) FOR I=1 TO N-4. FAC 200
C TWO RESULT VECTORS G,H ARE USED TO STORE THE ELEMENTS T(I,N-1) FOR FAC 210
C I=1 TO N-4 AND T(I,N) FOR I=1 TO N-3 RESPECTIVELY. THE ELEMENTS FAC 220
C T(N-3,N-1), T(N-2,N-1), T(N-2,N) AND T(N-1,N) ARE THEN GIVEN BY FAC 230
C A(N-3)+G(N-3), B(N-2)+G(N-2), A(N-2)+H(N-2), B(N-1)+H(N-1) FAC 240
C RESPECTIVELY. FAC 250
  DIMENSION A(N), B(N), C(N), G(N), H(N) FAC 260
  N1 = N - 1 FAC 270
  N2 = N - 2 FAC 280
  N3 = N - 3 FAC 290
  N4 = N - 4 FAC 300
  C(1) = SQRT(C(1)) FAC 310
  V = B(1)/C(1) FAC 320
  C(2) = SQRT(C(2)-V*V) FAC 330
  B(1) = V/C(2) FAC 340
  DO 10 I=3,N2 FAC 350
    I1 = I - 1 FAC 360
    I2 = I - 2 FAC 370
    U = A(I2)/C(I2) FAC 380
    V = B(I1)/C(I1) - B(I2)*U FAC 390
    C(I) = SQRT(C(I)-U*U-V*V) FAC 400
    A(I2) = U/C(I) FAC 410
    B(I1) = V/C(I) FAC 420
  10 CONTINUE FAC 430
  G(1) = A(N1)/C(1) FAC 440
  G(2) = -B(1)*G(1) FAC 450
  U = A(N3)/C(N3) FAC 460
  V = B(N2)/C(N2) - B(N3)*U FAC 470
  DO 20 I=3,N2 FAC 480
    G(I) = -B(I-1)*G(I-1) - A(I-2)*G(I-2) FAC 490
  20 CONTINUE FAC 500
  W = 0.0 FAC 510
  DO 30 I=1,N4 FAC 520
    W = W + G(I)*G(I) FAC 530
  30 CONTINUE FAC 540
  C(N1) = SQRT(C(N1)-W-(G(N3)+U)**2-(G(N2)+V)**2) FAC 550
  W = 1.0/C(N1) FAC 560
  A(N3) = U*W FAC 570
  B(N2) = V*W FAC 580
  DO 40 I=1,N2 FAC 590
    G(I) = G(I)*W FAC 600
  40 CONTINUE FAC 610
  H(1) = B(N)/C(1) FAC 620
  H(2) = A(N)/C(2) - B(1)*H(1) FAC 630
  U = A(N2)/C(N2) FAC 640
  V = B(N1)/C(N1) - B(N2)*U FAC 650
  DO 50 I=3,N1 FAC 660
    H(I) = -B(I-1)*H(I-1) - A(I-2)*H(I-2) FAC 670
  50 CONTINUE FAC 680
  W = 0.0 FAC 690
  DO 60 I=1,N2 FAC 700
    W = W + G(I)*H(I) FAC 710
  60 CONTINUE FAC 720
  H(N1) = H(N1) - W - G(N2)*U FAC 730
  W = 0.0 FAC 740
  DO 70 I=1,N3 FAC 750
    W = W + H(I)*H(I) FAC 760
  70 CONTINUE FAC 770
  C(N) = SQRT(C(N)-W-(H(N2)+U)**2-(H(N1)+V)**2) FAC 780
  W = 1.0/C(N) FAC 790
  A(N2) = U*W FAC 800
  B(N1) = V*W FAC 810
  DO 80 I=1,N1 FAC 820
    H(I) = H(I)*W FAC 830
  80 CONTINUE FAC 840
  RETURN FAC 850
  END FAC 860

SUBROUTINE SOLVE(E, F, G, H, Q, N) SOL 10
C THE SUBROUTINE SOLVES THE SET OF N LINEAR EQUATIONS RTZ=Q WHERE T IS SOL 20
C AN UPPER TRIANGULAR MATRIX WITH UNIT ELEMENTS ON THE DIAGONAL AND R SOL 30
C DENOTES THE TRANSPOSE OF T. THE OTHER NON-ZERO ELEMENTS OF T OCCUR ON SOL 40
C THE SUPER-DIAGONAL I.E. T(I,I+1), THE SUPER-SUPER-DIAGONAL I.E. SOL 50

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ALGORITHM 513

Analysis of In-Situ Transposition [F1]

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Key Words and Phrases: transposition in place, matrix transposition, permutation
CR Categories: 4.49, 5.9, 5.30
Language: Fortran

DESCRIPTION

Algorithms for in-place matrix transposition make use of the cyclic structure of the transposition mapping. A clear introduction to the process was given in 1958 by Windley [6]. In 1968 Boothroyd presented ACM Algorithm 302 [1]. Then in 1970 Laffin and Brefner presented ACM Algorithm 380 which they showed to be significantly faster than Algorithm 302 [4].

In this paper a number of theoretical results describing the transposition process are derived. Algorithms 302 and 380 are described and applications of some of the theoretical results to accelerate both algorithms are described. Finally, experimental comparison of the performance of Algorithm 380 and the improved version is presented.

Theoretical Development

Fortran conventions are assumed. The $m \times n$ matrix A is stored by column in a vector \vec{y} of adjacent storage locations. It is convenient to number components of \vec{y} from 0 to $mn - 1$. Using this convention, matrix element $A(I, J)$ is stored in y_k , where $k = m(J - 1) + I - 1$. The transposition process moves the contents of y_k to y_p where $p = n(I - 1) + J - 1$. Thus the transposition process can be regarded as a permutation mapping acting on the indices of \vec{y} . The symbol π is used to denote this mapping, and the set of indices which π permutes is denoted by S . With this notation, $S = [0, 1, \dots, mn - 1]$ and the formal equation $\pi(k) = p$ indicates that the contents of storage location y_k is moved to location y_p in the transposition process.

The theoretical material is organized in three sections: additional notation, development of general properties of π , and description of the fixed point properties of π .

Notation

Symbols used in this paper are as follows.

- $C(x)$ Cycle of π containing element $x \in S$. $C(x)$ is generated by repeated application of π to x , i.e. $C(x) = [x, \pi(x), \pi^2(x), \dots]$.
 $L(x)$ Length of the cycle $C(x)$.

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- $P(d)$ Number of cycles of length d .
 $\mu(n)$ Möbius function; for n having prime factorization, $n = P_1^{e_1} \dots P_k^{e_k}$,

$$\mu(n) = \begin{cases} 1 & \text{if } n = 1, \\ (-1)^k & \text{if } e_1 = e_2 = \dots = e_k = 1, \\ 0 & \text{otherwise.} \end{cases}$$

 $[a, b]$ Greatest common divisor of a and b .
 $d | u$ d is a divisor of u .

General Properties of π

An analytic expression for π is given in [1]. That is, for $x \in S$,

$$\pi(x) = \begin{cases} nx \bmod mn - 1 & \text{if } x \neq mn - 1, \\ mn - 1 & \text{if } x = mn - 1. \end{cases} \quad (1)$$

In this paper the symbols x and y represent any element of S except $mn - 1$, and all algebraic expressions involving these are assumed evaluated mod $mn - 1$.

That π is a linear mapping is obtained directly from (1). That is,

$$(ax + by) = a\pi(x) + b\pi(y). \quad (2)$$

This property is used to relate π to its fixed points.

The inverse mapping is used in both Algorithms 302 and 380 to traverse the cycles in the reverse direction.

THEOREM 1. $\pi^{-1}(y) = \begin{cases} my & \text{for } y \in S, y \neq mn - 1, \\ mn - 1 & \text{for } y = mn - 1. \end{cases}$

PROOF. $\pi^{-1}(\pi(y)) = \pi(\pi^{-1}(y)) = mny \bmod mn - 1 = ((mn - 1)y + y) \bmod mn - 1 = y$.

The transposition algorithms depend upon the cyclic structure of π . Two theorems are presented which describe the properties of these cycles. Theorem 2 shows how to find the length of the longest cycle (i.e. the cycles starting with $x = 1$, $A(2, 1)$). It also demonstrates that the length of any cycle must be a divisor of the length of the longest cycle. Theorem 3 provides a number theoretic formula which gives the number of cycles of any length. No application has been found for Theorem 3, but it may be of interest to other researchers. The interested reader is referred to [2] for the proof.

THEOREM 2. *The order of π is $L(1)$.*

PROOF. The order of a permutation mapping is the least common multiple of the lengths of its cycles [5]. Let $d = L(1)$. Then $n^d = 1 \bmod mn - 1$. It is sufficient to show, for any $y \in S$, that $\pi^d(y) = y$. If $y = mn - 1$, then the theorem is trivially true. For $y \neq mn - 1$, $\pi^d(y) = n^d y \bmod mn - 1 = y$.

THEOREM 3. *Assume that $k > 1$. The number of cycles of length k , $\frac{1}{k}P(k)$, is given by $P(k) = k^{-1} \sum_{d|k} \mu(k/d)[n^d - 1, mn - 1]$.*

The number of fixed points ($k = 1$) is discussed in the next section.

Fixed Point Properties of π

It is mainly through the symmetry with respect to the fixed points that computational leverage is gained. The mapping always has at least two fixed points, 0 and $mn - 1$. If S has a midpoint, then the midpoint is also a fixed point. Formally,

THEOREM 4. *If m and n are odd, then $\frac{1}{2}(mn - 1)$ is a fixed point.*

PROOF. Let $n = 2k + 1$. By eq. (1), $\pi[\frac{1}{2}(mn - 1)] = \frac{1}{2}(2k + 1)(mn - 1) = \frac{1}{2}(mn - 1)$

Two formulas for calculation of the number of fixed points of π are now given. Let n_f represent the number of fixed points of π .

THEOREM 5. $n_f - 1 = [n - 1, mn - 1] = [m - 1, mn - 1]$, and the fixed points are uniformly spaced at $0, d, 2d, \dots$, where $d = (mn - 1)/(n_f - 1)$.

PROOF. Because of eq. (1), all the fixed points of π except $mn - 1$ satisfy the congruence $x(n - 1) = 0 \bmod mn - 1$. The solutions of this congruence are $0, d, 2d, \dots$, where $d = (mn - 1)/[n - 1, mn - 1]$. (See [5, ch. 8].) Clearly π and π^{-1} have the same fixed points; so $n_f - 1 = [m - 1, mn - 1]$ by symmetry.

THEOREM 6. $n_f - 1 = [m - 1, n - 1]$.

PROOF. Let $d = [m - 1, n - 1]$. By Theorem 5, $(n_f - 1) \mid d$. Conversely, the identity $(n - 1)(m + 1) = mn - 1 + ((n - 1) - (m - 1))$ shows that $d \mid mn - 1$. This and Theorem 5 show that $d \mid (n_f - 1)$. Hence $d = n_f - 1$.

The mapping π is symmetric with respect to fixed points in the sense that for fixed point a , where $a \neq mn - 1$,

$$\pi(a \pm x) = a \pm \pi(x). \quad (3)$$

The fixed point 0 is usually the only fixed point satisfying eq. (3). That is, if $m \neq 1$ and $n \neq 1$, the probability that π has exactly two fixed points (0 and $mn - 1$) is the probability that $m - 1$ and $n - 1$ are relatively prime (Theorem 6). This probability is $6/\pi^2 = .608$ [3, ch. 18]. The symmetry with respect to 0 is applied in Algorithm 380 and extended in this paper.

When $a = 0$, the nontrivial part of eq. (3) is

$$\pi(-x) = -\pi(x). \quad (4)$$

It is evident from eq. (4) that $C(-x) = -C(x)$, and consequently that $L(x) = L(-x)$. Equation (4) was proved by Laffin and Brefner [4] by more complicated methods and was the basis for their improved algorithm. We call the symmetric cycle $C(-x)$ the companion cycle. The cycle and the companion cycle in some cases coincide. In this event the length of the cycle is an even integer, and x and $-x$ are separated by half the cycle. That is,

THEOREM 7. $C(x) = C(-x)$ if and only if $-x = n^k x$, where $k = L(x)/2$.

PROOF. Since $-x \in C(x)$ there is a smallest integer k such that $-x = n^k x$. This implies that $x = n^k(-x)$, which in turn implies that $x = n^k(n^k x)$. Thus $L(x) \mid 2k$. Since $k < L(x)$, we must have $L(x) = 2k$. Conversely, if $n^k x = -x$, then $-x \in C(x)$ and $C(x) = C(-x)$.

Algorithm Improvement

Algorithms 302 and 380 are now described. For each algorithm, efficiency improvements based on the theoretical development in the previous section are discussed. Finally, timing comparisons are given.

Algorithm 302 [1]. For $i = 1, \dots, mn - 1$, compute the index $j = \pi^{-1}(i)$ and exchange y_i and y_j , provided $j \geq i$. The case $j < i$ indicates that the element originally in y_j is now elsewhere following previous exchanges. Then search the cycle $C(j)$ for the first index $j_r \geq i$ and exchange y_{j_r} and y_i . It terminates when y_{mn-2} is filled.

Symmetry with respect to the fixed point 0 [eq. (4)] can be exploited to accelerate Algorithm 302. The idea is to process the array from both ends toward the middle simultaneously. That is, process element i and $-i \bmod (mn - 1)$ concurrently. This reduces the time spent searching cycles. The revised algorithm is then:

1. Set $i = 0$.
2. Increment i . If $i > (mn/2 - 1)$, terminate. Otherwise, perform step 3 ($mn/2$ is truncated to integer value because of Theorem 4).
3. Compute j as follows: Let r be the smallest positive integer such that $i \leq \pi^{-r}(i) \leq mn - 1 - i$. Then $j = \pi^{-r}(i)$.
4. If $j = i$, return to step 2. Otherwise, exchange y_i and y_j and perform step 5.
5. If $j = mn - 1 - i$, return to step 2. Otherwise, perform the symmetric exchange of y_{mn-1-j} and y_{mn-1-i} . (If $j = mn - 1 - i$, this exchange was already performed in step 4.) Return to step 2.

Algorithm 380 [4]. First the number of fixed points is determined by testing each element in the array to see if $\pi(i) = i$. Then the array is scanned from lowest to highest subscript. Each array subscript in turn is checked to see if that subscript is the smallest subscript in its cycle. Whenever such an element is found, the entire cycle is moved. If the companion cycle is distinct, then the companion cycle is moved after movement of the cycle is complete. The algorithm keeps track of the total number of elements moved and terminates when all (mn) elements have

been moved. It is interesting to note that Algorithm 380 works its way backwards around each cycle, using the inverse, saving itself a fetch and store for each element moved. Algorithm 380 can achieve greater efficiency if scratch storage is available. The scratch storage is utilized to keep track of the elements which have already been moved.

The improvements to Algorithm 380 are based again on the symmetry with respect to 0 and upon the calculations of the number of fixed points. Theorem 6 allows use of Euclid's algorithm to calculate the number of fixed points. Theorem 7 and eq. (4) allow the processing of a cycle and its companion concurrently. When the cycle and its companion coincide, the process is terminated halfway through the cycle (Theorem 7). We also used the linearity to calculate $\pi(x+1) = \pi(x) + n$ to start the next cycle. It is interesting to note that on the computer used (CDC 6600 using the RUN compiler), the calculation of $mx \bmod mn - 1$ using the system modulo functions is slower than using the formulation used in Algorithm 380.

Timing Comparisons

The timing tests utilized a set of 21 test matrices which were presented in [4]. The per-element transposition time is defined to be the time required to transpose a matrix divided by the number of elements in the matrix. Each algorithm was tested on the complete set of test matrices and the average per-element transposi-

Table I

Algorithm	Average per-element transposition time (μ s)
Algorithm 302	92.2
Algorithm 302 (revised)	67.1
Algorithm 380 (without scratch storage)	70.7
Algorithm 380 (with $(m+n)/2$ scratch storage)	46.9
Algorithm 380 (revised) (without scratch storage)	49.2
Algorithm 380 (revised) (with $(m+n)/2$ scratch storage)	28.5

Table II

Size $M \times N$	IWRK = $(M+N)/2$			IWRK = 1		
	Original T1*	Revised T2*	$\frac{2 \times (T1 - T2)}{(T1 + T2)}$	Original T3*	Revised T4*	$\frac{2 \times (T3 - T4)}{(T3 + T4)}$
7 \times 60	1.58	.84	.61	2.45	1.63	.40
7 \times 70	1.71	.86	.66	2.52	1.62	.43
7 \times 80	3.09	2.01	.42	4.15	2.96	.33
7 \times 90	2.32	1.21	.63	5.01	3.52	.35
7 \times 100	2.50	1.28	.65	3.88	2.46	.45
8 \times 60	1.63	.84	.64	1.66	.87	.62
8 \times 70	3.46	2.29	.41	4.94	3.6	.31
8 \times 80	4.50	3.10	.37	6.69	5.02	.29
8 \times 90	2.53	1.25	.68	2.55	1.22	.71
8 \times 100	3.96	2.41	.49	7.28	5.13	.35
9 \times 60	2.44	1.42	.53	4.12	2.85	.36
9 \times 70	3.32	2.11	.45	5.65	4.02	.34
9 \times 80	2.56	1.22	.71	2.49	1.24	.67
9 \times 90	2.81	1.45	.64	4.13	2.66	.43
9 \times 100	3.28	1.84	.56	7.44	5.43	.31
45 \times 50	11.19	7.07	.45	19.39	14.20	.31
45 \times 60	9.54	4.70	.68	9.22	4.66	.66
46 \times 50	16.19	11.23	.36	25.49	19.57	.26
46 \times 60	17.49	12.37	.34	27.96	21.57	.26
47 \times 50	18.45	13.54	.31	25.90	20.37	.24
47 \times 60	9.29	4.94	.61	9.65	4.89	.65

* Times are given in seconds/100.

tion time was calculated. These are summarized in Table I. A detailed comparison of the best algorithms (380 and the revised 380) is presented in Table II in the format utilized in [4]. All tests were performed using a CDC 6600, with compilation performed by the RUN compiler.

The simple modification of Algorithm 302 resulted in about a 25-percent efficiency improvement (Table I). The performance improvement of the revised Algorithm 380 varied from 20 percent to 70 percent on individual test matrices (Table II) with an average of about 35 percent.

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ALGORITHM

SUBROUTINE TRANS(A, M, N, MN, MOVE, IWRK, IOK)	TRA	10
C *****	TRA	20
C ALGORITHM 380 - REVISED	TRA	30
C *****	TRA	40
C A IS A ONE-DIMENSIONAL ARRAY OF LENGTH MN=M*N, WHICH	TRA	50
C CONTAINS THE MXN MATRIX TO BE TRANSPOSED (STORED	TRA	60
C COLUMNWISE). MOVE IS A ONE-DIMENSIONAL ARRAY OF LENGTH IWRK	TRA	70
C USED TO STORE INFORMATION TO SPEED UP THE PROCESS. THE	TRA	80
C VALUE IWRK=(M+N)/2 IS RECOMMENDED. IOK INDICATES THE	TRA	90
C SUCCESS OR FAILURE OF THE ROUTINE.	TRA	100
C NORMAL RETURN IOK=0	TRA	110
C ERRORS IOK=-1 ,MN NOT EQUAL TO M*N	TRA	120
C IOK=-2 ,IWRK NEGATIVE OR ZERO	TRA	130
C IOK.GT.0, (SHOULD NEVER OCCUR),IN THIS CASE	TRA	140
C WE SET IOK EQUAL TO THE FINAL VALUE OF I WHEN THE SEARCH	TRA	150
C IS COMPLETED BUT SOME LOOPS HAVE NOT BEEN MOVED	TRA	160
C NOTE * MOVE(I) WILL STAY ZERO FOR FIXED POINTS	TRA	170
DIMENSION A(MN), MOVE(IWRK)	TRA	180
C CHECK ARGUMENTS AND INITIALIZE.	TRA	190
IF (M.LT.2 .OR. N.LT.2) GO TO 120	TRA	200
IF (MN.NE.M*N) GO TO 180	TRA	210
IF (IWRK.LT.1) GO TO 190	TRA	220
IF (M.EQ.N) GO TO 130	TRA	230
NCOUNT = 2	TRA	240
K = MN - 1	TRA	250
DO 10 I=1,IWRK	TRA	260
MOVE(I) = 0	TRA	270
10 CONTINUE	TRA	280
IF (M.LT.3 .OR. N.LT.3) GO TO 30	TRA	290
C CALCULATE THE NUMBER OF FIXED POINTS, EUCLIDS ALGORITHM	TRA	300
C FOR GCD(M-1,N-1).	TRA	310
IR2 = M - 1	TRA	320
IR1 = N - 1	TRA	330
20 IR0 = MOD(IR2,IR1)	TRA	340
IR2 = IR1	TRA	350
IR1 = IR0	TRA	360
IF (IR0.NE.0) GO TO 20	TRA	370
NCOUNT = NCOUNT + IR2 - 1	TRA	380
C SET INITIAL VALUES FOR SEARCH	TRA	390
30 I = 1	TRA	400
IM = M	TRA	410
C AT LEAST ONE LOOP MUST BE RE-ARRANGED	TRA	420
GO TO 80	TRA	430
C SEARCH FOR LOOPS TO REARRANGE	TRA	440
40 MAX = K - I	TRA	450
I = I + 1	TRA	460
IF (I.GT.MAX) GO TO 160	TRA	470
IM = IM + M	TRA	480
IF (IM.GT.K) IM = IM - K	TRA	490
I2 = IM	TRA	500

IF (I.EQ.I2) GO TO 40	TRA 510
IF (I.GT.IWRK) GO TO 60	TRA 520
IF (MOVE(I).EQ.0) GO TO 80	TRA 530
GO TO 40	TRA 540
50 I2 = M*I1 - K*(I1/N)	TRA 550
60 IF (I2.LE.I .OR. I2.GE.MAX) GO TO 70	TRA 560
I1 = I2	TRA 570
GO TO 50	TRA 580
70 IF (I2.NE.I) GO TO 40	TRA 590
C REARRANGE THE ELEMENTS OF A LOOP AND ITS COMPANION LOOP	TRA 600
80 I1 = I	TRA 610
KMI = K - I	TRA 620
B = A(I1+1)	TRA 630
I1C = KMI	TRA 640
C = A(I1C+1)	TRA 650
90 I2 = M*I1 - K*(I1/N)	TRA 660
I2C = K - I2	TRA 670
IF (I1.LE.IWRK) MOVE(I1) = 2	TRA 680
IF (I1C.LE.IWRK) MOVE(I1C) = 2	TRA 690
NCOUNT = NCOUNT + 2	TRA 700
IF (I2.EQ.I) GO TO 110	TRA 710
IF (I2.EQ.KMI) GO TO 100	TRA 720
A(I1+1) = A(I2+1)	TRA 730
A(I1C+1) = A(I2C+1)	TRA 740
I1 = I2	TRA 750
I1C = I2C	TRA 760
GO TO 90	TRA 770
C FINAL STORE AND TEST FOR FINISHED	TRA 780
100 D = B	TRA 790
B = C	TRA 800
C = D	TRA 810
110 A(I1+1) = B	TRA 820
A(I1C+1) = C	TRA 830
IF (NCOUNT.LT.MN) GO TO 40	TRA 840
C NORMAL RETURN	TRA 850
120 IOK = 0	TRA 860
RETURN	TRA 870
C IF MATRIX IS SQUARE, EXCHANGE ELEMENTS A(I, J) AND A(J, I).	TRA 880
130 N1 = N - 1	TRA 890
DO 150 I=1, N1	TRA 900
J1 = I + 1	TRA 910
DO 140 J=J1, N	TRA 920
I1 = I + (J-1)*N	TRA 930
I2 = J + (I-1)*M	TRA 940
B = A(I1)	TRA 950
A(I1) = A(I2)	TRA 960
A(I2) = B	TRA 970
140 CONTINUE	TRA 980
150 CONTINUE	TRA 990
GO TO 120	TRA 1000
C ERROR RETURNS.	TRA 1010
160 IOK = I	TRA 1020
170 RETURN	TRA 1030
180 IOK = -1	TRA 1040
GO TO 170	TRA 1050
190 IOK = -2	TRA 1060
GO TO 170	TRA 1070
END	TRA 1080

REMARK ON ALGORITHM 513

Analysis of In-Situ Transposition [F1]

[E.G. Cate and D.W. Twigg, *ACM Trans. Math. Software* 3, 1 (March 1977), 104-110]

and

REMARK ON ALGORITHM 467

Matrix Transposition in Place [F1]

[N. Brenner, *Comm. ACM* 16, 11 (Nov. 1973), 692-694]

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It is somewhat distressing to note the publication of Algorithm 513. While the algorithm is, as claimed, an improvement over Algorithm 380 [1], it is by no means as good as Algorithm 467, which was published long before Algorithm 513 was even submitted for publication. A reanalysis of the timing figures in Algorithms 467 and 513 suggests very strongly that Algorithm 467 is superior to Algorithm 513 except, possibly, when the modulus of the generating function for the transposition permutation is prime.

A direct comparison of the algorithms was performed using the Fortran G (nonoptimized) compiler on an IBM 370/158 running under VM/CMS. The results, reported in Table I, indicate quite clearly that Algorithm 467 enjoys a nearly 2:1 advantage except when the modulus is prime. It should be noted that it is not so much that Algorithm 467 is degraded in this case as that Algorithm 513 performs exceptionally well.

Table I. Timing Results for Algorithms 467 and 513
(All times are in milliseconds. Work area sizes are $(R + C)/2$ in all cases.)

Rows	Columns	Modulus	467	513	513/467
45	50	13.173	58	88	1.52
45	60	2699	72	68	.94
45	180	7.13.89	212	444	2.08
45	200	8999	214	199	.93
46	50	11 ² .19	58	128	2.21
46	60	31.89	72	147	2.04
46	180	17.487	213	359	1.69
46	200	9199	212	181	.85
47	50	34.29	54	134	2.46
47	60	2819	66	62	.94
47	180	11.769	207	411	1.99
47	200	2.13.241	430	628	1.46

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ALGORITHM 514

A New Method of Cubic Curve Fitting Using Local Data [E2]

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Key Words and Phrases: interpolation, cubic splines, local data
CR Categories: 5.13
Language: Algol

DESCRIPTION

This algorithm is a local interpolation method to find an approximating function $y = f(x)$ through a set of given data points (x_i, y_i) with $x_1 < \dots < x_i < x_{i+1} < \dots$. Probably the best known methods for this are cubic splines [3, 6, 7] and the "classical" Lagrangian polynomial interpolation, of which Aitken's method is perhaps the most stable implementation [2]. Both of these methods are nonlocal, i.e. the interpolated values depend on all the data points. Under many circumstances, however, it is desirable to use a local interpolation method for which the values of the interpolating function depend only on a few of the nearest data points. For example, a local method may be preferred in an interactive system where one does not want to recalculate a complete picture when one point is changed, or where one wants to save storage in multidimensional interpolation.

The method used in our algorithm is similar to cubic spline interpolation in that different cubic, or third degree, polynomials $y = f_i(x)$ represent the function between different pairs of data points $(x_i, y_i) - (x_{i+1}, y_{i+1})$ and in that the coefficients are chosen to ensure continuity of the function and its first derivative at each data point. We have, however, sacrificed the spline's continuity of the second derivative in order to ensure that the coefficients depend only on the local data points.

A similar approach is described by Maude [4] which also ensures the continuity of the function and its first derivative at each data point. However, Maude's functions, at their simplest, reduce to quintic (fifth degree) polynomials, whereas ours are cubic polynomials with resulting computational advantages. Another method with a similar goal is due to Akima [1]. But whereas Akima's method is designed from geometrical considerations, ours is algebraic and appears to have some practical advantages. Thus our method gives an exact fit to the original polynomial $f(x)$ if the data points satisfy a cubic equation $y_i = f(x_i)$. As a result, in general it is more accurate for data of mathematical origin, or where smooth curves are expected. Our algorithm also appears to give more spline-like functions than Akima's, and in particular tends to produce much smaller discontinuities in

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the second derivatives at the data points. This is particularly important when interpolating surfaces, $z = f(x, y)$, by interpolating successively in two independent directions, x and y ; for when depicting a perspective view of the surface, for example, a discontinuity in the second derivative of the function often results in a noticeable irregularity (a kink) in a visible edge.

Tables I and II compare the performance of our algorithm with other interpolation methods, and give the results obtained from a program run on a Control Data 7600 which applies six methods to data generated from five mathematical functions. Table I shows the average deviation found between the interpolated result and the true value from the original function, while Table II gives the average discontinuity of the second derivative at the nodes. The data used for the interpolations were generated for the following 18 unequally spaced values of x : $-2.95, -2.6, -2.1, -1.8, -1.4, -1.0, -0.75, -0.3, -0.05, 0.2, 0.55, 0.9, 1.25, 1.6, 1.7, 2.1, 2.4, 3.0$.

The first three methods are nonlocal ones: Aitken's method and two cubic spline methods. The first of the two cubic splines (a) is as implemented in Späth [7] and is probably the most used of the spline routines in the NAG Library [5], while the second cubic spline (b) was implemented by Reid [6]. They differ mainly in their treatment of the end intervals: Späth sets the second derivative at the endpoints to zero; Reid forces the third derivative at the points adjacent to the ends to be continuous. The second three methods given in Table I are local ones: Akima's, Maude's, and the method described below.

Our method is implemented in the two procedures described here. For each point (x_i, y_i) the first procedure, *find gradients at nodes*, initially finds that cubic poly-

Table I. Average Deviation from Original Function for Various Interpolation Methods ($\times 10^6$)

Function	Method					
	Nonlocal			Local		
	Aitken's	Cubic splines (a) (b)		Akima's	Maude's	Present
x^3	0	28051	0	25709	8172	0
x^4	0	165921	4561	112169	68566	6434
$e^{-x^2/2}$	36	164	76	1260	640	78
$\tanh x$	10409	140	147	833	619	214
$\sin x$	0*	297	116	1531	1055	176

* The actual value was $0.011 (\times 10^{-6})$. Aitken's method gives a remarkably accurate fit for $\sin x$ and a remarkably poor one for $\tanh x$.

Table II. Average Discontinuity of Second Derivative at Data Points ($\times 10^4$)

Function	Method			
	All non-local methods as in Table I	Local		
		Akima's	Maude's	Present
x^3	0	79953	0	0
x^4	0	391953	0	5766
$e^{-x^2/2}$	0	5027	0	268
$\tanh x$	0	4844	0	469
$\sin x$	0	6293	0	198

nomial which passes through the three points (x_{i-1}, y_{i-1}) , (x_i, y_i) , (x_{i+1}, y_{i+1}) and which gives a least squares fit to the two neighboring data points (x_{i-2}, y_{i-2}) , (x_{i+2}, y_{i+2}) . As written, in the least squares fitting the weights given to the last two points are not equal, but are inversely proportional to the squares of their distances from x_{i-1} and x_{i+1} , respectively; we have found experimentally that such weighting leads to smaller discontinuities in the second derivatives at the nodes. The gradient of this cubic at the central point (x_i, y_i) is then calculated. This process is repeated for all the points (x_i, y_i) , except for the first two and the last two data points for which the subscripts x_{i-2} etc. or x_{i+2} etc. would be out of range; for these extreme points a cubic is fitted exactly through the first four or the last four points as appropriate.

The second procedure, *coeffs*, then constructs as the approximating function for the interval $x_i \leq x \leq x_{i+1}$ that cubic polynomial which has the correct ordinates y_i and y_{i+1} and the calculated gradients at the two endpoints. The procedure calculates the four coefficients c_0, c_1, c_2, c_3 for which the approximating polynomial may be expressed in the form

$$f(x) = c_0 + c_1(x - x_i) + c_2(x - x_i)^2 + c_3(x - x_i)^3.$$

This is preferred to the form $a_0 + a_1x + a_2x^2 + a_3x^3$ which is more unstable against roundoff errors.

Normally the procedure *find gradients at nodes* need only be called once for any given set of x_i, y_i , thus avoiding unnecessary repetition of calculation; the gradients it calculates are then used by the procedure *coeffs* for any required interval.

We have also found the method useful in two-dimensional interpolation to fit a surface $z = f(x, y)$ through data on a rectangular grid (x_i, y_j, z_{ij}) . This may be achieved by first applying the method to all the rows of data, keeping y_j fixed and varying i , and then by applying it again to the resulting four sets of coefficients, but this time keeping x_i fixed and varying j .

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We would like to thank M.J.D. Powell and the referees for a number of valuable suggestions. We also would like to acknowledge the assistance provided by Mrs. S. Costello in the computation of the various figures shown in the tables.

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ALGORITHM

```

procedure find gradients at nodes (x,y,grad,lbx,ubx);
  value lbx,ubx; integer lbx,ubx;
  array x,y,grad;
comment x and y are arrays with subscripts from lbx
  to ubx giving the abscissae and ordinates respectively
  of the data points. These should be in
  ascending order of abscissae, with no two abscissae

```

equal. The arrays must contain at least four points.

grad is an array having the same subscripts as x and y into which will be placed the calculated gradients of the required cubics at each of the data points.

lbx and ubx are integers giving the minimum and maximum values respectively of the subscripts of the arrays x, y, and grad;

```

begin
  integer i,iless2,iless1,iplus1,iplus2;
  real x0,x1,x2,x3,x4,y2,prod1,prod2,num,denom,g,
    coeff2,xdiff,xprod,weight;
  for i:= lbx step 1 until ubx do
    begin
      comment special treatment is needed at end points;
      illess1:= if i > lbx then i-1 else i+3;
      iplus1:= if i < ubx then i+1 else i-3;
      x2:= x[i]; y2:= y[i];
      x1:= x[iless1]-x2; x3:= x[iplus1]-x2;
      comment first fit a quadratic through x1,x2,x3;
      prod1:= (y[iless1]-y2)*x3;
      prod2:= (y[iplus1]-y2)*x1;
      denom:= x1*x3*(x[iless1]-x[iplus1]);
      g:= (x1*prod2-x3*prod1)/denom;
      coeff2:= (prod1-prod2)/denom;
      comment if x0 exists, find its contribution to the
      cubic adjustment;
      if i ≤ lbx+1 then num:= denom:= 0.0 else
        begin
          illess2:= i-2; x0:= x[iless2]-x2;
          xdiff:= x[iless2]-x[iless1];
          xprod:= x0*xdiff*(x[iless2]-x[iplus1]);
          weight:= xprod/(xdiff*xdiff);
          num:= weight*(y[iless2]-y2-x0*(g+x0*coeff2));
          denom:= weight*xprod
        end;
      comment if x4 exists, find its contribution to the
      cubic adjustment;
      if i < ubx-1 then
        begin
          iplus2:= i+2; x4:= x[iplus2]-x2;
          xdiff:= x[iplus2]-x[iplus1];
          xprod:= x4*xdiff*(x[iplus2]-x[iless1]);
          weight:= xprod/(xdiff*xdiff);
          num:= num+weight*(y[iplus2]-y2-x4*(g+x4*coeff2));
          denom:= denom+weight*xprod
        end;
      grad[i]:= g+num*x1*x3/denom
    end
  end of procedure find gradients at nodes;

  procedure coeffs(x,y,grad,i,c0,c1,c2,c3);
  value i; integer i;
  real c0,c1,c2,c3;
  array x,y,grad;

```

comment This procedure calculates the coefficients of the cubic which has values $y[i]$, $y[i+1]$ and gradients $\text{grad}[i]$, $\text{grad}[i+1]$ at $x[i]$, $x[i+1]$ respectively.

The cubic takes the form

$$c0 + c1(x-x[i]) + c2(x-x[i])^2 + c3(x-x[i])^3.$$

To interpolate as suggested in the Description, above, these coefficients would be used for values of x in the range $x[i] \leq x \leq x[i+1]$;

begin

real h, dy ;

$c0 := y[i]$;

$h := x[i+1] - x[i]$; $dy := y[i+1] - c0$;

$c1 := \text{grad}[i]$;

$c2 := (3.0 \times dy - h \times (2.0 \times c1 + \text{grad}[i+1])) / (h \times h)$;

$c3 := (h \times (c1 + \text{grad}[i+1]) - 2.0 \times dy) / h^3$

end of procedure coeffs;

ALGORITHM 515

Generation of a Vector from the Lexicographical Index [G6]

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Key Words and Phrases: combinations

CR Categories: 5.39

Language: Fortran

DESCRIPTION

Let $C = \{C_1, C_2, \dots, C_m\}$ be the set of combinations of n items taken p at a time and arranged in lexicographical order. Given an integer, i ($1 \leq i \leq m$, $n \geq p > 0$), the algorithm derives C_i . Previous algorithms [1, 3, 6] have accomplished this by the generation of all vectors between an initial point in the combination space and the desired vector. The cited algorithms are computationally advantageous if all combinations or a sequential subset are required. However, the algorithm given here is advantageous if a few randomly selected combinations are needed and each selection is not based on the previous selection history.

A one-to-one correspondence between the universe of n items and the first n natural numbers is established. The combinations produced by the algorithm are selections of p of the first n integers in lexicographical order [2, 5]. Thus the combinations produced by the algorithm may be regarded equivalently as pointers to combinations of any objects. If the combinations are generated sequentially (according to the index), exactly the same order is achieved as that of Mifsud's algorithm [4].

Vector generation is performed by probing the combination space and sequentially reducing the interval of uncertainty within which the indexed vector resides. Beginning with the leftmost vector position, trial values in ascending order for the digits are tested by computing the index of the (partial) trial vector, using essentially the method of Walter [7]. The lexicographical index, i , for a combination C_i is computed as follows: Let the elements of the combination vector be denoted X_1, X_2, \dots, X_p . Using $X_0 = 0$, the index i is

$$i = 1 + \sum_{j=1}^p Q_j; \quad Q_j = \begin{cases} \sum_{k=X_{j-1}+1}^{X_j-1} \binom{n-k}{p-j} & \text{if } X_{j-1} + 1 \leq X_j - 1 \\ 0 & \text{otherwise.} \end{cases}$$

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This formula is the inverse of the present algorithm. The technique consists of accumulating terms of the form of those in the expression for Q_j until a total greater than or equal to the desired vector index is reached. The value causing this condition is X_j . The sum is reduced by the amount of the last term (so that Q_j only contains terms through $X_j - 1$) and the search resumes at the next position to the right. The probing halts when there are no positions remaining. A previous ACM algorithm [8] is used to compute the binomial coefficients.

Once a value has been eliminated from a position, it is eliminated from all following positions since the values are produced in ascending order. Thus the number of probes required to find a vector is exactly the numeric value of the last element. That is, the vector {1, 3, 7} will be found in seven probes given any combination space containing seven or more things taken three at a time. This leads to an expression for the average number of probes required to find a vector based entirely on the frequency of values occurring as the last element:

$$\binom{n}{p}^{-1} \sum_{m=0}^{n-p} \binom{n-1-m}{p-1} (n-m).$$

It can be seen that each nonzero term in Q_p (i.e. Q_j for $j = p$) equals unity. If the first term (1) in the expression for i is combined with Q_p , it can be written in the form

$$Q_p = X_j - X_{j-1}.$$

Taking advantage of this identity, an improvement in performance may be obtained by computing

$$C_p = C_{p-1} + L - K,$$

where L and K are variables given within the algorithm.

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ALGORITHM

SUBROUTINE COMB(N, P, L, C)	COM	10
C THIS SUBROUTINE FINDS THE COMBINATION SET OF N THINGS	COM	20
C TAKEN P AT A TIME FOR A GIVEN LEXICOGRAPHICAL INDEX.	COM	30
C N - NUMBER OF THINGS IN THE SET	COM	40
C P - NUMBER OF THINGS IN EACH COMBINATION	COM	50
C L - LEXICOGRAPHICAL INDEX OF COMBINATION SOUGHT	COM	60
C C - OUTPUT ARRAY CONTAINING THE COMBINATION SET	COM	70

C THE FOLLOWING RELATIONSHIPS MUST EXIST AMONG THE INPUT	COM 80
C VARIABLES. L MUST BE GREATER THAN OR EQUAL TO 1 AND LESS	COM 90
C THAN OR EQUAL TO THE MAXIMUM LEXICOGRAPHICAL INDEX.	COM 100
C P MUST BE LESS THAN OR EQUAL TO N AND GREATER THAN ZERO.	COM 110
INTEGER N, P, L, C(P), K, R, P1, BINOM	COM 120
C INITIALIZE LOWER BOUND INDEX AT ZERO	COM 130
K = 0	COM 140
C LOOP TO SELECT ELEMENTS IN ASCENDING ORDER	COM 150
P1 = P - 1	COM 160
DO 20 I=1,P1	COM 170
C SET LOWER BOUND ELEMENT NUMBER FOR NEXT ELEMENT VALUE	COM 180
C(I) = 0	COM 190
IF (I.NE.1) C(I) = C(I-1)	COM 200
C LOOP TO CHECK VALIDITY OF EACH ELEMENT VALUE	COM 210
10 C(I) = C(I) + 1	COM 220
R = BINOM(N-C(I),P-I)	COM 230
K = K + R	COM 240
IF (K.LT.L) GO TO 10	COM 250
K = K - R	COM 260
20 CONTINUE	COM 270
C(P) = C(P1) + L - K	COM 280
RETURN	COM 290
END	COM 300
INTEGER FUNCTION BINOM(M, N)	BIN 10
C ACM ALGORITHM 160 TRANSLATED TO FORTRAN. CALCULATES THE	BIN 20
C NUMBER OF COMBINATIONS OF M THINGS TAKEN N AT A TIME.	BIN 30
INTEGER M, N, P, I, N1, R	BIN 40
N1 = N	BIN 50
P = M - N1	BIN 60
IF (N1.GE.P) GO TO 10	BIN 70
P = N1	BIN 80
N1 = M - P	BIN 90
10 R = N1 + 1	BIN 100
IF (P.EQ.0) R = 1	BIN 110
IF (P.LT.2) GO TO 30	BIN 120
DO 20 I=2,P	BIN 130
R = (R*(N1+I))/I	BIN 140
20 CONTINUE	BIN 150
30 BINOM = R	BIN 160
RETURN	BIN 170
END	BIN 180

ALGORITHM 516

An Algorithm for Obtaining Confidence Intervals and Point Estimates Based on Ranks in the Two-Sample Location Problem [G1]

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Pennsylvania State University

Key Words and Phrases: asymptotic linearity, confidence interval, Illinois method, point estimate, ranks, regula falsi, robust, Wilcoxon-Mann-Whitney rank test
CR Categories: 5.5
Language: Fortran

DESCRIPTION

Suppose we are sampling from two populations which have the same continuous distribution except for a possible shift in location. For notational convenience let X_1, \dots, X_m be a random sample from $F(x)$ and let Y_1, \dots, Y_n be a random sample from $F(x - \theta)$. In testing hypotheses concerning θ the robust two-sample rank test (Wilcoxon-Mann-Whitney) is often used, but the related point estimate and confidence intervals are used relatively rarely. The reasons for this are primarily computational. While it has been practical to carry out the test for large sample sizes, it has not been practical for the estimates.

The most common method for computing the estimates is based on the $m \cdot n$ ordered differences $D_{ij} = Y_j - X_i$ (for instance, see Noether [5, p. 138]). These differences can be quickly obtained and ordered for small m and n ; however, for large m and n they may be costly to obtain, even with the use of a computer. The storage is high (for example, if m and n are 1000, then 1,000,000 words of storage are needed); furthermore, computation time is high (of order $m \cdot n$, even if the samples are ordered in a preliminary step).

One way of avoiding the computation difficulty is to modify the estimate, for instance, as Antille [1] did in one sample problem. We propose instead an iteration procedure based on the function $U(\theta) = \#(Y_j - X_i \leq \theta)$. This is the Mann-Whitney test statistic for testing that the shift is θ ; it differs only by a constant from the Wilcoxon rank sum statistic. It is an increasing function of θ which steps up one unit at each difference $Y_j - X_i$. The point estimate and confidence interval can be found by essentially inverting this function; for example, the point estimate of θ is $\hat{\theta} = \text{median}(D_{ij})$, which can be found by solving the equation $U(\theta) = mn/2$ (the middle such value if it is not unique). The endpoints of the

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confidence interval can be obtained similarly. See Hettmansperger and McKean [3] for a graphical view of $U(\theta)$.

The algorithm given here finds the point estimate and confidence interval by solving such equations. Since $U(\theta)$ is monotone, methods such as regula falsi (linear interpolation) may be used. Of the methods tried, we have found the Illinois method, a modification of the regula falsi method outlined by Dowell and Jarratt [2], to be the best. The number of iterations is reduced by a factor of 2 or 3 when it is compared to regula falsi; hence we have incorporated it into the algorithm. The algorithm requires a preliminary and separate ordering of the X and Y . We have found that this ordering allows an algorithm for calculating $U(\theta)$ which saves a considerable amount of time.

The algorithm eliminates the storage problem. Furthermore, we feel the algorithm obtains the estimates quickly enough to make them more attractive. The time for an evaluation of $U(\theta)$ is linear in sample size but the number of iterations decreases due to the asymptotic linearity of $U(\theta)$ (see Jureckova [4]). For example, in a test run with $m = 1000$ and $n = 1200$ the computation time was 0.19 CPU seconds to obtain both the point estimate and confidence interval on an IBM 370 Model 168 (this costs about \$0.02).

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ALGORITHM

```

SUBROUTINE RANKCI(X, M, Y, N, PERC, DPOINT, DLOW, DHIGH,          RAN 10
* IERR, ILOW, IHIGH)                                           RAN 20
C SUBROUTINE TO FIND THE POINT ESTIMATE AND CONFIDENCE INTERVAL RAN 30
C RELATED TO THE TWO-SAMPLE RANK TEST (MANN-WHITNEY, WILCOXON), RAN 40
C FOR THE PARAMETER D = MU(Y) - MU(X), WHERE MU IS A MEASURE OF RAN 50
C LOCATION.                                                    RAN 60
C THE STATISTICAL ASSUMPTIONS ARE THAT X AND Y ARE RANDOM SAMPLES RAN 70
C TAKEN INDEPENDENTLY FROM TWO DIFFERENT POPULATIONS, AND THAT RAN 80
C THE POPULATIONS HAVE THE SAME DISTRIBUTION EXCEPT FOR LOCATION. RAN 90
C IN PARTICULAR, IT IS ASSUMED THAT THE SCALE PARAMETERS (VARIANCES, RAN 100
C IF THEY EXIST) ARE EQUAL.                                     RAN 110
C THIS ROUTINE FINDS THE POINT ESTIMATE FOR D AS THAT VALUE DPOINT RAN 120
C SUCH THAT THE MANN-WHITNEY U STATISTIC ACHIEVES ITS EXPECTED RAN 130
C VALUE FOR THE TEST OF THE NULL HYPOTHESIS D = DPOINT (THE MIDDLE RAN 140
C SUCH VALUE IF IT IS NOT UNIQUE).                             RAN 150
C THE CONFIDENCE INTERVAL IS FOUND AS THOSE VALUES OF DD SUCH THAT RAN 160
C THE NULL HYPOTHESIS D = DD IS NOT REJECTED BY THE TWO SAMPLE RANK RAN 170
C TEST.                                                         RAN 180
C THE METHOD OF CALCULATING THE TEST STATISTIC U IS THE ORIGINAL RAN 190
C METHOD OF MANN AND WHITNEY, MODIFIED TO TAKE ADVANTAGE OF HAVING RAN 200
C THE DATA IN ORDER. THIS SPEEDS CALCULATIONS CONSIDERABLY, WHICH RAN 210
C IS IMPORTANT SINCE THE METHODS USED MUST EVALUATE THE STATISTIC RAN 220
C REPEATEDLY.                                                 RAN 230
C THE POINT AND CONFIDENCE INTERVALS ARE FOUND BY ITERATION, USING RAN 240
C TRIMMED MEANS AS STARTING POINTS. THE BASIC ITERATION PROCEDURE RAN 250
C IS A MODIFICATION OF THE REGULA FALSI (LINEAR INTERPOLATION) RAN 260
C METHOD, (DOWELL + JARRATT, 1971,BIT) WHICH CONVERGES QUICKLY RAN 270
C BECAUSE OF THE ASYMTOTIC LINEARITY OF THE U STATISTIC AS A RAN 280
C FUNCTION OF D (JURECKOVA, 1971, ANN. MATH. STAT.) RAN 290
C THE PERCENTAGE OF THE CONFIDENCE INTERVAL IS OBTAINED BY A NORMAL RAN 300
C APPROXIMATION, USING CONTINUITY CORRECTION. RAN 310
C INPUT IS AS FOLLOWS RAN 320
C X - ARRAY OF DATA OF DIMENSION M IN NON-DECREASING ORDER RAN 330
C M - NUMBER OF OBSERVATIONS IN X SAMPLE (AT LEAST 2) RAN 340
C Y - ARRAY OF DATA OF DIMENSION N IN NON-DECREASING ORDER RAN 350

```



```

10 UMED1 = XM*XN/2.0 - 0.5          RAN 1120
   UMED2 = UMED1 + 1.0             RAN 1130
   IODD = 0                       RAN 1140
C                                     P IS ONE-TAILED PROB.          RAN 1150
20 P = (100.0-PERC)/200.0         RAN 1160
C                                     FIND MEAN AND VARIANCE OF THE U-STATISTIC RAN 1170
C                                     FOR USE IN A NORMAL APPROXIMATION. RAN 1180
   UMU = XM*XN/2.0                RAN 1190
   UVAR = XM*XN*(XM+XN+1.0)/12.0  RAN 1200
   USIGMA = SQRT(UVAR)            RAN 1210
C                                     NORMAL APPROX. (WITH CONTINUITY CORRECTION) IS RAN 1220
C                                     P = PHI ( (ULOW+.5-UMU)/USIGMA ) RAN 1230
C                                     (WHERE PHI IS STANDARD NORMAL DIST. FUNCTION, RAN 1240
C                                     AND PHINV IS ITS INVERSE). RAN 1250
   ULOW = USIGMA*PHINV(P) + UMU - 0.5 RAN 1260
C                                     ROUND CRITICAL VALUE DOWN TO INTEGER RAN 1270
C                                     (THIS GIVES CONSERVATIVE BOUNDS), AND FIND P. RAN 1280
   IU = ULOW                      RAN 1290
   IF (IU.LT.0) IU = 0            RAN 1300
   ULOW = IU                      RAN 1310
   Z = (ULOW+0.5-UMU)/USIGMA     RAN 1320
   P = PHI(Z)                    RAN 1330
   PERC = 100.0*(1.0-2.0*P)     RAN 1340
C                                     WANT TO INVERT FUNCTION U AT A HALF INTEGER RAN 1350
   ULOW = ULOW + 0.5            RAN 1360
C                                     UHIGH RAN 1370
   UHIGH = XM*XN - ULOW         RAN 1380
   ILOW = ULOW + .5            RAN 1390
   IHIGH = UHIGH + .5          RAN 1400
C                                     ILOW AND IHIGH GIVE THE ORDERED DIFFERENCES RAN 1410
C                                     WHICH FORM THE CONFIDENCE INTERVAL. RAN 1420
C                                     AN ESTIMATE OF THE SLOPE OF THE LINEAR RAN 1430
C                                     APPROXIMATION TO FMANN. RAN 1440
   SLOPE = (ABS(Z)*SQRT(XM*XN*(XM+XN)))/((DH-DL)*SQRT(3.)) RAN 1450
C                                     X1 AND X2 BRACKET THE LOWER CONFIDENCE LIMIT. RAN 1460
C                                     THEN BY CALLING THE ROUTINE ILL THE LOWER CONFIDENCE RAN 1470
C                                     LIMIT WITHIN EPS2 IS RETURNED VIA DLOW. RAN 1480
   CALL BRACK(DL, ULOW, SLOPE, X, M, Y, N, X1, FX1, X2, FX2) RAN 1490
   CALL ILL(DLOW, ULOW, X1, FX1, X2, FX2, X, M, Y, N, EPS2) RAN 1500
C                                     SAME FOR UPPER END VIA DHIGH. RAN 1510
   CALL BRACK(DH, UHIGH, SLOPE, X, M, Y, N, X1, FX1, X2, FX2) RAN 1520
   CALL ILL(DHIGH, UHIGH, X1, FX1, X2, FX2, X, M, Y, N, EPS2) RAN 1530
C                                     A NEW ESTIMATE OF SLOPE BASED ON THE CONFIDENCE RAN 1540
C                                     INTERVAL (DLOW,DHIGH), UNLESS THE LENGTH OF RAN 1550
C                                     THE INTERVAL IS SMALLER THAN EPS2. RAN 1560
   IF (DHIGH.GT.DLOW+EPS2) SLOPE = ((DH-DL)/(DHIGH-DLOW))*SLOPE RAN 1570
C                                     THE SAME ROUTINES ARE USED FOR THE ESTIMATE RAN 1580
C                                     DPOINT EXCEPT EPS1 IS USED. THE MIDPOINT RAN 1590
C                                     OF THE CONFIDENCE INTERVAL WILL BE THE INITIAL RAN 1600
C                                     ESTIMATE OF DPOINT. RAN 1610
   D = (DLOW+DHIGH)/2.0         RAN 1620
   CALL BRACK(D, UMED1, SLOPE, X, M, Y, N, X1, FX1, X2, FX2) RAN 1630
   CALL ILL(DPOINT, UMED1, X1, FX1, X2, FX2, X, M, Y, N, EPS1) RAN 1640
C                                     IF M*N IS ODD THE ESTIMATE IS DPOINT, RAN 1650
   IF (IODD.EQ.1) RETURN       RAN 1660
C                                     IF EVEN THEN THE VALUE DPOINT IS THE LOWER RAN 1670
C                                     CENTER ESTIMATE. THE UPPER CENTER ESTIMATE IS D2, RAN 1680
C                                     AND THE FINAL ESTIMATE IS THE AVERAGE OF THE TWO. RAN 1690
   CALL BRACK(DPOINT, UMED2, SLOPE, X, M, Y, N, X1, FX1, X2, FX2) RAN 1700
   CALL ILL(D2, UMED2, X1, FX1, X2, FX2, X, M, Y, N, EPS1) RAN 1710
   DPOINT = (DPOINT+D2)/2.     RAN 1720
   RETURN                       RAN 1730
   END                           RAN 1740

SUBROUTINE CHECK(X, M, Y, N, PERC, IERR) CHE 10
C SUBROUTINE TO DO ERROR CHECKING FOR RANKCI CHE 20
C WRITTEN JUNE 1975 BY T. RYAN CHE 30
  DIMENSION X(M), Y(N) CHE 40
C                                     CHECK FOR INSUFFICIENT DATA CHE 50
  IF (M.GE.2 .AND. N.GE.2) GO TO 10 CHE 60
  IERR = 1 CHE 70
  RETURN CHE 80
C                                     CHECK FOR PROPER PERCENT CONFIDENCE CHE 90
10 IF (PERC.LT.99.999 .AND. PERC.GT.49.999) GO TO 20 CHE 100
  IERR = 2 CHE 110
  RETURN CHE 120

```


C	CHECK FOR X AND Y ARRAYS IN NON-DECREASING ORDER	CHE	130
	20 DO 30 I=2,M	CHE	140
	IF (X(I-1).GT.X(I)) GO TO 50	CHE	150
	30 CONTINUE	CHE	160
	DO 40 I=2,N	CHE	170
	IF (Y(I-1).GT.Y(I)) GO TO 50	CHE	180
	40 CONTINUE	CHE	190
	RETURN	CHE	200
C	X OR Y ARRAY OUT OF ORDER	CHE	210
	50 IERR = 3	CHE	220
	RETURN	CHE	230
	END	CHE	240
	SUBROUTINE TMEAN(Z, N, ZBAR, VARZB)	TME	10
C	GIVEN DATA Z(1), Z(2), ..., Z(N) IN NON-DECREASING ORDER,	TME	20
C	THIS ROUTINE FINDS THE ALPHA-TRIMMED MEAN ZBAR, AND THE	TME	30
C	ESTIMATED VARIANCE OF ZBAR, VARZB.	TME	40
C	WRITTEN JUNE 1975 BY T. RYAN	TME	50
	DIMENSION Z(N)	TME	60
	DATA ALPHA /0.10/	TME	70
	XN = N	TME	80
C	NT IS NUMBER TRIMMED FROM EACH END	TME	90
C	N1 IS NUMBER OF LOWEST OBSERVATION NOT TRIMMED	TME	100
C	N2 IS NUMBER OF HIGHEST OBSERVATION NOT TRIMMED	TME	110
	NT = ALPHA*XN	TME	120
	N1 = NT + 1	TME	130
	N2 = N - NT	TME	140
C	TRIMMED MEAN	TME	150
	SUM = 0.0	TME	160
	DO 10 I=N1,N2	TME	170
	SUM = SUM + Z(I)	TME	180
	10 CONTINUE	TME	190
	X = N - 2*NT	TME	200
	ZBAR = SUM/X	TME	210
C	WINSORIZED SUM OF SQUARES	TME	220
	SUM = 0.0	TME	230
	DO 20 I=N1,N2	TME	240
	SUM = SUM + (Z(I)-ZBAR)**2	TME	250
	20 CONTINUE	TME	260
	IF (NT.EQ.0) GO TO 30	TME	270
	XNT = NT	TME	280
	SUM = SUM + XNT*(Z(N1-1)-ZBAR)**2 + XNT*(Z(N2+1)-ZBAR)**2	TME	290
	30 VARZB = SUM/(XN*XN)	TME	300
	RETURN	TME	310
	END	TME	320
	SUBROUTINE ILL(XVAL, FVAL, X1, F1, X2, F2, X, M, Y, N, EPS)	ILL	10
C	THIS ROUTINE SOLVES THE EQUATION	ILL	20
C	F(T)-FVAL=0	ILL	30
C	FOR A MONOTONE FUNCTION F, IN THIS INSTANCE FMANN. THE METHOD	ILL	40
C	USED IS THE ILLINOIS METHOD AS DESCRIBED BY DOWELL AND JARRATT	ILL	50
C	(1971,BIT), EXCEPT FOR A MODIFICATION WHEN CLOSE TO THE ROOT.	ILL	60
C	THEN IF THE ROOT WAS NOT TRAPPED ON THE LAST ITERATION THE ROU-	ILL	70
C	TINE SWITCHES TO THE BISECTION METHOD.	ILL	80
C	INPUT -	ILL	90
	FVAL = THE VALUE OF THE FUNCTION AT THE ROOT.	ILL	100
	X1 AND X2 ARE VALUES WHICH BRACKET THE ROOT, THAT IS	ILL	110
	EITHER	ILL	120
	F(X1) .LT. FVAL .LT. F(X2)	ILL	130
	OR	ILL	140
	F(X2) .LT. FVAL .LT. F(X1)	ILL	150
	FX1 = F(X1)	ILL	160
	FX2 = F(X2)	ILL	170
	X, M, Y, AND N ARE QUANTITIES USED BY THE FUNCTION FMANN	ILL	180
	EPS = THE ACCURACY OF THE SOLUTION.	ILL	190
	X1, X2, AND THEIR FUNCTIONAL VALUES ARE CHANGED THROUGHOUT THE	ILL	200
	ROUTINE.	ILL	210
	OUTPUT -	ILL	220
	XVAL = THE ROOT WITHIN EPS.	ILL	230
	DIMENSION X(M), Y(N)	ILL	240
	F1 = F1 - FVAL	ILL	250
	F2 = F2 - FVAL	ILL	260
	IBISEC = 0	ILL	270
	10 CONTINUE	ILL	280
	IF (ABS(X2-X1).LT.EPS) GO TO 40	ILL	290

```

C          X3 IS THE INTERSECTION OF THE SECANT LINE          ILL 300
C          FORMED BY (X1,F1), (X2,F2) AND THE X-AXIS.        ILL 310
X3 = X2 - (F2*(X2-X1))/(F2-F1)                               ILL 320
IF (IBISEC.EQ.1) X3 = (X1+X2)/2.                             ILL 330
IBISEC = 0                                                    ILL 340
CALL FMANN(X3, X, M, Y, N, F3)                                ILL 350
F3 = F3 - FVAL                                                ILL 360
IF (F3*F2) 20, 20, 30                                         ILL 370
C          ROOT WAS TRAPPED, SO USE REGULA FALSI            ILL 380
20 X1 = X2                                                    ILL 390
F1 = F2                                                       ILL 400
X2 = X3                                                       ILL 410
F2 = F3                                                       ILL 420
GO TO 10                                                       ILL 430
C          ROOT WAS NOT TRAPPED, SO USE ILLINOIS MOD-        ILL 440
C          IFICATION.                                        ILL 450
30 X2 = X3                                                    ILL 460
F2 = F3                                                       ILL 470
F1 = F1/2.                                                    ILL 480
IF (ABS(F2).LE.ABS(F1)) GO TO 10                               ILL 490
C          IF ILLINOIS MODIFICATION IS MORE RADICAL          ILL 500
C          THAN BISECTION, THEN USE BISECTION.              ILL 510
F1 = 2.0*F1                                                  ILL 520
IBISEC = 1                                                    ILL 530
GO TO 10                                                       ILL 540
40 XVAL = (X1+X2)/2.                                         ILL 550
RETURN                                                         ILL 560
END                                                           ILL 570

SUBROUTINE BRACK(XINIT, FVAL, SLOPE, X, M, Y, N, X1, FX1, X2,
* FX2)
C SUPPOSE THE FUNCTION F IS MONOTONE AND IT IS DESIRED TO SOLVE
C THE EQUATION
C          F(T) - FVAL = 0.
C THIS ROUTINE RETURNS TWO VALUES WHICH BRACKET THE ROOT.
C INPUT -
C          XINIT = INITIAL ESTIMATE OF THE ROOT.
C          FVAL = THE VALUE OF THE FUNCTION AT THE ROOT.
C          SLOPE = APPROXIMATE SLOPE OF THE FUNCTION IN A
C          NEIGHBORHOOD OF THE ROOT.
C          X, M, Y, AND N ARE QUANTITIES USED BY THE FUNCTION WHICH
C          IN THIS INSTANCE IS FMANN.
C NONE OF THE ABOVE QUANTITIES IS CHANGED THROUGHOUT THIS ROUTINE
C OUTPUT -
C          X1 AND X2 BRACKET THE ROOT, THAT IS EITHER
C          F(X1) .LT. FVAL .LT. F(X2)
C          OR
C          F(X2) .LT. FVAL .LT. F(X1)
C          FX1 = F(X1)
C          FX2 = F(X2)
C DIMENSION X(M), Y(N)
X1 = XINIT
CALL FMANN(X1, X, M, Y, N, FX1)
DELTA = 1.5*((FVAL-FX1)/SLOPE)
10 X2 = X1 + DELTA
CALL FMANN(X2, X, M, Y, N, FX2)
IF ((FX1-FVAL)*(FX2-FVAL).LT.0.) RETURN
X1 = X2
GO TO 10
END

SUBROUTINE FMANN(D, X, MM, Y, NN, U)
C ROUTINE TO CALCULATE THE MANN-WHITNEY STATISTIC U.
C U DIFFERS ONLY BY A CONSTANT FROM THE WILCOXON 2-SAMPLE RANK
C STATISTIC W.
C INPUT -
C          D = NULL HYPOTHESIS VALUE OF THE SHIFT OF THE Y
C          POPULATION FROM THE X POPULATION (I.E. THE NULL
C          HYPOTHESIS BEING TESTED IS MU(Y) = MU(X) + D).
C          X = ARRAY CONTAINING THE SAMPLE FROM ONE POPULATION,
C          WHICH MUST BE IN NON-DECREASING ORDER.
C          M = DIMENSION (SAMPLE SIZE) OF X
C          Y = ARRAY CONTAINING THE SAMPLE FROM THE OTHER POPULATION,
C          WHICH MUST BE IN NON-DECREASING ORDER.
C          N = DIMENSION (SAMPLE SIZE) OF Y

```

```

C NONE OF THE ABOVE IS ALTERED BY THIS ROUTINE. FMA 150
C OUTPUT - FMA 160
C U = MANN-WHITNEY TEST STATISTIC. FMA 170
C THIS ROUTINE IS INTENDED TO BE USED IN AN ITERATION ROUTINE FOR FMA 180
C ESTIMATION ONLY, SINCE NO CHECKS OR ADJUSTMENTS ARE MADE FOR TIES. FMA 190
C THIS ROUTINE IS DESIGNED TO BE VERY FAST, SINCE IT IS TO BE USED FMA 200
C IN AN ITERATION PROCEDURE, SO NO CHECKS ARE MADE TO INSURE THAT FMA 210
C M AND N ARE POSITIVE, OR THAT X AND Y ARE NON-DECREASING. FMA 220
C ALL CHECKING SHOULD BE DONE IN CALLING PROGRAM. FMA 230
C U= SUM (NO. OF Y(J) LESS THAN (OR EQUAL) TO X(I)+DELTA) WHERE FMA 240
C THE SUM IS OVER I. FMA 250
C SINCE ARRAYS X AND Y ARE ORDERED, IF X(I) IS GREATER THAN Y(J), FMA 260
C X(I+1) MUST ALSO BE. FMA 270
C WRITTEN 3/75 BY T. RYAN. LAST UPDATED 6/75 BY T. RYAN FMA 280
C DIMENSION X(MM), Y(NN) FMA 290
C DELTA = D FMA 300
C M = MM FMA 310
C N = NN FMA 320
C JLE IS THE NUMBER OF Y VALUES LESS THAN OR FMA 330
C EQUAL TO X(I). FMA 340
C IU IS ACCUMULATED AS THE VALUE OF THE U FMA 350
C STATISTIC. FMA 360
C JLE = 0 FMA 370
C IU = 0 FMA 380
C MAIN LOOP FMA 390
C DO 30 I=1,M FMA 400
C XI = X(I) + DELTA FMA 410
10 IF (XI.LT.Y(JLE+1)) GO TO 20 FMA 420
C JLE = JLE + 1 FMA 430
C IF (JLE.GE.N) GO TO 50 FMA 440
C GO TO 10 FMA 450
20 IU = IU + JLE FMA 460
30 CONTINUE FMA 470
40 U = IU FMA 480
C RETURN FMA 490
C X(I) IS GREATER THAN (OR EQUAL) TO ALL Y(J). FMA 500
C THEREFORE X(I+1), ..., X(M) ARE ALL GREATER FMA 510
C THAN ALL Y(J). FMA 520
50 IU = IU + (M-I+1)*N FMA 530
C GO TO 40 FMA 540
C END FMA 550

FUNCTION PHINV(P) PHI 10
C INVERSE NORMAL DISTRIBUTION FUNCTION. PHI 20
C IF Z IS N(0,1), THIS FUNCTION RETURNS PHINV DEFINED BY PHI 30
C  $P(A .LT. PHINV) = P, (0.0 .LT. P \text{ AND } P .LT. 1.0)$ . PHI 40
C REF. HANDBOOK OF MATHEMATICAL FUNCTIONS, 1964, USDC, NATIONAL PHI 50
C BUREAU OF STANDARDS, WASH. DC. P. 933, FORMULA 26.2.23. PHI 60
C ERROR WILL BE LESS THAN  $4.5E-04$ . PHI 70
C WRITTEN 4/76 BY T. RYAN BASED ON ROUTINE BY H. D. KNOBLE (1966). PHI 80
C IF (P.LE.0. .OR. P.GE.1.0) GO TO 30 PHI 90
C IF (P.LT.0.5) GO TO 10 PHI 100
C SIGN = 1.0 PHI 110
C PTAIL = 1.0 - P PHI 120
C GO TO 20 PHI 130
10 SIGN = -1.0 PHI 140
C PTAIL = P PHI 150
20 T = SQRT(ALOG(1.0/(PTAIL*PTAIL))) PHI 160
C Z = ABS(T-(2.515517+T*(0.802853+T*0.010328))/(1.0+T* PHI 170
C * (1.432788+T*(0.189269+T*0.001308)))) PHI 180
C PHINV = Z*SIGN PHI 190
C RETURN PHI 200
30 CONTINUE PHI 210
C RETURN PHI 220
C END PHI 230

FUNCTION PHI(X) PHI 10
C NORMAL DISTRIBUTION FUNCTION PHI 20
C LET Z BE N(0,1), THEN PHI IS DEFINED BY  $\text{PHI} = P(Z .LE. X)$ . PHI 30
C REF. HANDBOOK OF MATHEMATICAL FUNCTIONS, 1964, USDC, NATIONAL PHI 40
C BUREAU OF STANDARDS, WASH. DC, P. 933, FORMULA 26.2.19. PHI 50
C ERROR IS LESS THAN  $1.5E-07$ . PHI 60
C WRITTEN 4/76 BY T. RYAN, BASED ON ROUTINE BY H. D. KNOBLE (1966). PHI 70
C Z = X PHI 80
C IF (Z.LT.0.0) GO TO 10 PHI 90

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COLLECTED ALGORITHMS (cont.)

516-P 8- 0

ISIGN = 1	PHI 100
GO TO 20	PHI 110
10 Z = -X	PHI 120
ISIGN = -1	PHI 130
20 P = 0.5*(1.0+Z*(0.049867347+Z*(0.0211410061+Z*(3.2776263D-3+Z*	PHI 140
* (3.80036E-05+Z*(4.88906E-05+Z*5.383E-06))))))**(-16)	PHI 150
IF (ISIGN.EQ.1) P = 1.0 - P	PHI 160
PHI = P	PHI 170
RETURN	PHI 180
END	PHI 190

ALGORITHM 517

A Program for Computing the Condition Numbers of Matrix Eigenvalues Without Computing Eigenvectors [F2]

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Key Words and Phrases: eigenvalue, condition number
CR Categories: 5.14
Language: Fortran

DESCRIPTION

1. Theoretical Background

1.1 *The Sensitivity of Eigenvalues.* Several good programs are available for the computation of the eigenvalues of real and complex matrices [2, 3, 7]. Because of the limitations of finite precision arithmetic, these programs cannot produce, in general, the exact eigenvalues of the given matrix A . However the computed numbers are always (very close to) the eigenvalues of a matrix $A + E$ which is very close to A . This matrix E is not unique and error analyses [6] have shown the existence of E 's with satisfactorily small upper bounds on $\|E\| / \|A\|$. Here $\|\cdot\|$ denotes an appropriate matrix norm.

It follows from these remarks that a good program will not always deliver accurate approximations to the eigenvalues of A . It can happen that some, or all, of the eigenvalues are very sensitive to changes in the matrix elements; so some, or all, of the eigenvalues of $A + E$ may differ sharply from those of A . Actually this is true only for non-normal matrices. Real symmetric matrices—indeed all normal matrices—determine their eigenvalues very well; the change induced in an eigenvalue of such an A cannot exceed the spectral norm of E (which is defined below).

Two questions arise: How can this sensitivity be measured, and how cheaply can it be computed?

Simple eigenvalues. To any simple eigenvalue λ of A there correspond both a column vector x and a row eigenvector y^* (the conjugate transpose of y) which are unique to within a scalar multiple. Thus

$$Ax = x\lambda, \quad y^*A = \lambda y^* \quad (1)$$

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and $x = y$ if A is normal (i.e. $A^*A = AA^*$). The most popular measure of the sensitivity of λ was called the spectral condition number $\text{cond}(\lambda)$ by Wilkinson [6]. Let θ denote the acute angle between x and y ; then

$$\text{cond}(\lambda) \equiv \secant \theta = \|y\| \cdot \|x\| / |y^*x|$$

where

$$\|v\| = \sqrt{(v^*v)}. \quad (2)$$

This definition gives a number in $[1, \infty)$ which increases monotonically with the sensitivity of λ to changes in A .

In order to justify this definition two popular matrix norms will be used:

$$\begin{aligned} \|M\| &\equiv \max_{v \neq 0} \|Mv\| / \|v\| = \sqrt{(\lambda_{\max}(M^*M))}, \\ \|M\|_E &\equiv \sqrt{(\sum_i \sum_j |m_{ij}|^2)} = \sqrt{(\text{trace}(M^*M))}. \end{aligned} \quad (3)$$

Let $|\delta\lambda|$ be the change in λ corresponding to a change δA in A . It can be shown that

$$\text{cond}(\lambda) = \sup |\delta\lambda| / \|\delta A\|_E \quad (4)$$

over all non-null infinitesimal δA .

Another useful characterization of $\text{cond}(\lambda)$ is the following: The spectral projector P_λ of λ is the matrix which projects every vector into a multiple of the eigenvector of λ . It is easy to verify that for simple λ

$$P_\lambda = xy^*/y^*x \quad (5)$$

where y^*x is a scalar, and by using the fact that P_λ is of rank 1, one can show that

$$\text{cond}(\lambda) = \|P_\lambda\|_E = \|P_\lambda\|. \quad (6)$$

It is this characterization which can be generalized.

Multiple eigenvalues. When λ has geometric multiplicity m , almost all perturbations of A break λ into m simple eigenvalues in such a way that $\sup |\delta\lambda| / \|\delta A\|$ is unbounded. Thus it is customary to set $\text{cond}(\lambda) = \infty$ in this case.

There is more to be said however. A reasonable definition (see Kahan [4]) puts

$$\text{cond}(\lambda) \equiv \sup |\delta\lambda| / \|\delta A\| \quad (7)$$

over all non-null infinitesimal δA which preserve the multiplicity of λ . This number can be estimated because $\text{cond}(\lambda) \leq \|P_\lambda\|_E / m$ where the spectral projector P_λ satisfies $AP_\lambda = P_\lambda A = \lambda P_\lambda + N_\lambda$ and N_λ is nilpotent (i.e. $N_\lambda^m = 0$). Moreover P_λ can be found from the expression

$$P_\lambda = X(Y^*X)^{-1}Y^* \quad (8)$$

where the columns of X and rows of Y^* are bases for the invariant subspaces of λ .

We have followed the usual practice (of setting the condition number equal to infinity) in our program CONDIR but wish to point out that it is feasible to bring into adjacent positions on the diagonal of the Schur form any associated ill-conditioned eigenvalues. The spectral projector for this group of eigenvalues can then be found from (8), and if its norm is small, the group can be designated as a cluster. That is another project.

If more specific information is required, then the individual elements of P_λ will be involved because

$$\partial\lambda/\partial a_{ij} = e_j^* P_\lambda e_i \quad (9)$$

with λ simple.

A warning should be offered at this point. The measures presented above are based on the Euclidean vector norm and the convention that A acts on vectors in Euclidean n -space. It can happen that this model is quite inappropriate for certain applications and then the conventional condition numbers will be irrelevant. However it is only the order of magnitude (base 10) of $\text{cond}(\lambda)$ which is wanted, in most cases, and this will be constant over a large range of norms.

1.2 *Invariance Properties.* When the role of the matrix is to be stressed the condition number is written $\text{cond}(\lambda, A)$.

THEOREM. *If Q is unitary, i.e. $Q^*Q = QQ^* = I$, and λ is a simple eigenvalue of A , then*

$$\text{cond}(\lambda, QAQ^*) = \text{cond}(\lambda, A).$$

PROOF. Let $Ax = x\lambda$, $y^*A = \lambda y^*$. Then

$$(QAQ^*)(Qx) = (Qx)\lambda, \quad (y^*Q^*)(QAQ^*) = \lambda(y^*Q^*).$$

Because λ is simple, $y^*x \neq 0$ and

$$\begin{aligned} \text{cond}(\lambda, QAQ^*) &= \|y^*Q^*\| \cdot \|Qx\| / |(y^*Q^*)(Qx)|, \\ &= \|y^*\| \cdot \|x\| / |y^*x|, \\ &= \text{cond}(\lambda, A), \end{aligned}$$

because the Euclidean norm is unitarily invariant. \square

COROLLARY. *If a given matrix B is reduced to Hessenberg form H by unitary similarities (such as Householder transformations) and the QR algorithm is applied to H to produce, in the limit, a quasi-triangular matrix T , then*

$$\text{cond}(\lambda, B) = \text{cond}(\lambda, T).$$

1.3 *The Use and Cost of Condition Numbers.* A computed eigenvalue λ of a given matrix A is an exact eigenvalue of many matrices including some close to A . Let $A + E$ designate one of the closest matrices. Provided $(\|E\|_E / \|A\|_E)^2$ is negligible, the error in λ is bounded by $\text{cond}(\lambda) \|E\|_E$. Error analyses [6] give an upper bound β on $(\|E\|_E / \|A\|_E)$ when the Householder/QR method is used. It follows that $\log_{10}(|\lambda| / \beta \cdot \text{cond}(\lambda) \cdot \|A\|_E)$ gives the number of decimal digits in λ which are assuredly correct.

When no figures in λ can be relied on, then a warning tag should be attached to λ for most applications. Conversely, when an adequate number of figures are certified as correct in each eigenvalue of A , then the subsequent calculations are placed on a sounder footing.

These estimates of the number of correct figures have proved useful in the comparison of rival eigenvalue programs and in debugging big programs of which the eigenvalue calculations were merely a part.

A natural question at this stage is how much extra does it cost to compute $\text{cond}(\lambda)$ as well as λ ? The answer must depend on whether the user also computes x and/or y along with λ . We focus on real matrices and real arithmetic.

(1) If a complete Jordan factorization $A = X\Lambda Y^*$ ($Y^*X = I$) is computed then each $\text{cond}(\lambda_i)$ can be found from the definition $\|x_i\| \|y_i^*\| / |y_i^*x_i|$ at negligible extra cost in storage and time. No special program is needed and this case will not be considered further. Few dependable Jordan factorization routines are currently available.

(2) If a program is used which yields X and Λ but not Y^* , then it is necessary to compute the triangular factorization $L_x U_x$ and store it in an extra array. Then $\text{cond}(\lambda_i) = \|e_i^* X^{-1}\| \cdot \|X e_i\|$. To invert X costs n^3 basic operations whereas X and Λ may be found in approximately $7n^3$ operations using the double QR transformation.

No special program is needed. The time penalty is slight but the extra storage requirement is substantial. This case will not be discussed further.

(3) The eigenvalues λ of A may be found (EISPACK path, ELMHES, HQR) in under $4\frac{1}{2}n^3$ operations with no supplementary $n \times n$ storage arrays provided A can be overwritten. This is the most interesting case. No extra arrays are needed for the computation of $\text{cond}(\lambda_i)$, $i = 1, \dots, n$, but the multiplication count rises to approximately $7n^3$. See Section 1.4 for more details. The $O(n^2)$ terms bring down the ratio of running times; the increase is approximately 50 percent (± 15 percent).

Our method is easily described. The given matrix A is reduced to Hessenberg form H by *orthogonal* similarity transformations. Then H is transformed to quasi-triangular Schur form T by the double QR algorithm working on the whole of H and not just the remaining principal submatrices. None of the orthogonal transforming matrices is retained. Finally the column and row eigenvectors of T are found, for each λ , by back substitution and then discarded immediately after $\text{cond}(\lambda)$ has been calculated.

By Theorem 1, $\text{cond}(\lambda, T) = \text{cond}(\lambda, A)$.

For simplicity all condition numbers exceeding 10^{30} are recorded as 10^{30} .

The program uses only real arithmetic even if A has complex eigenvalues.

1.4 *Operation Counts.* In [5] Parlett and Wang point out that straightforward counts of multiplications and additions are unreliable indicators of running times. Nevertheless they are good to within a factor of 2 and they do give insight into the way the algorithm spends its time. An op is defined as a scalar multiplication or division followed by an addition.

ORTHES. The $(n - j)$ th step transforms the last j rows and columns while reducing column $(n - j)$ to upper Hessenberg form.

Row operations: $A \rightarrow A' = A - w\gamma(w^T A)$, $\gamma = 2/w^T w$, $w^T = (0, \dots, 0, x, \dots, x)$

Computation	γ	$w^T A$	$v^T = \gamma(w^T A)$	$A - wv^T$	Total
Cost	j	j^2	j	j^2	$2j(j + 1)$

Column operations: $A' \rightarrow A'' = A' - \gamma A' w w^T$

Computation	γ	$A' w$	$u = \gamma A' w$	$A' - u w^T$	Total
Cost	0	nj	n	nj	$n(2j + 1)$

Grand total: $\sum_{j=1}^{n-1} n(2j + 1) + 2j(j + 1) = \frac{5}{3}n(n^2 - 1)$

The program ELMHES is approximately twice as fast as ORTHES but will not preserve condition numbers.

CONDIT. It suffices to assume that all eigenvalues are real. To find the column and row eigenvectors for the j th eigenvalue requires backsolving triangular systems of $(j - 1)$ and $(n - j - 1)$ equations, respectively.

Computation	x	y^*	Cond
Cost	$\sum_{i=1}^{j-1} i$	$\sum_{i=1}^{n-j-1} i$	n

Grand total: $\frac{1}{3}n^3 + \frac{1}{2}n^2 + \frac{1}{6}n$

HQR. A typical double QR transformation acts on a $j \times j$ submatrix of a Hessenberg matrix. To restore column k to Hessenberg form requires the following operations for $k \leq j - 2$. For $k = j - 1$ the cost is $3j + 6$ ops.

Computation	Key quantities	Row operations	Column operations
Cost	9	$\sum_{l=k}^j 5$	$\sum_{l=1}^{\min(k+3, j)} 5$

Total: $\sum_{k=1}^{j-2} [9 + 5(j - k + 1) + 5(k + 3)] + 3j + 6 = 5j^2 + 22j - 54$

Assume four initial full transformations with $j = n$ and then two iterations per eigenvalue.

Grand total: $\frac{10}{3}n^3 + 47n^2 + \frac{11}{3}n - 182$

QR2N0Z. The same transformations as in HQR must act on the whole matrix. This changes the range of the row operation and not the column because the $j \times j$ submatrix being transformed is the leading principal submatrix.

Computation	Key quantities	Row operations	Column operations
Cost	9	$\sum_{l=k}^n 5$	$\sum_{l=1}^{\min(k+3, j)} 5$

$$\text{Total: } \sum_{k=1}^{j-2} [9 + 5(n - k + 1) + 5(k + 3)] + 3n + 21 = (5n + 29)(j - 2) + (3n + 21)$$

With the same assumptions as above,

$$\text{Grand total: } 5n^3 + 40n^2 + 23n - 300$$

A summary of operation counts is:

$$\begin{aligned} \text{ELMHES + HQR:} & \quad 4\frac{1}{2}n^3 + 47n^2 + 3n - 182; \\ \text{ORTHES + HQR:} & \quad 5n^3 + 47n^2 + 2n - 182; \\ \text{ORTHES + QR2NOZ + CONDIT:} & \quad 7n^3 + 40\frac{1}{2}n^2 + 21\frac{1}{2}n - 300. \end{aligned}$$

The actual timings were more favorable to our program than these operation counts suggest. The assumption of two iterations per eigenvalue is unrealistic. In practice there are more iterations with larger values of j and fewer with small values. With $20 \leq n \leq 60$ our program ran, on the average, 50 percent longer than did ELMHES + HQR; the worst case ran 65 percent longer.

2. Applicability

The program accepts real square matrices which can be stored in the high speed memory of the computer.

The condition numbers of all eigenvalues of all normal matrices (and this includes symmetric matrices) are unity and consequently the program is intended for use with non-normal matrices.

Before our programs QR2NOZ and CONDIT are used, A should be reduced to Hessenberg form H by orthogonal congruences. We recommend the procedure ORTHES described in [7] and its Fortran counterpart ORTHES described in [2, p. 297].

Our program QR2NOZ is an adaptation of HQR2 [2, p. 248] and is designed to avoid the formation of the product of all the similarity transformations used in the double QR algorithm and the calculation of the eigenvectors of the final matrix of the QR sequence.

3. Organizational Details

3.1 *Standardization.* (i) In the course of the QR algorithm applied to H it is possible for two real eigenvalues to be found at the same time as the roots of a 2×2 diagonal block

$$\begin{pmatrix} \alpha & \beta \\ \gamma & \delta \end{pmatrix}.$$

It is convenient in such cases to do a supplementary plane rotation which will reduce this diagonal block to upper triangular form and change the corresponding rows and columns of H accordingly.

If this transformation is done at the time the eigenvalues λ_1 and λ_2 are recorded then some of the quantities which determine the correct angle of rotation will be available.

This device is employed in HQR2 and has been carried over to QR2NOZ. The details are given below.

The parameters $c = \cos \theta$, $s = \sin \theta$ are determined so that

$$\begin{pmatrix} c & -s \\ s & c \end{pmatrix} \begin{pmatrix} \alpha & \beta \\ \gamma & \delta \end{pmatrix} \begin{pmatrix} c & s \\ -s & c \end{pmatrix}$$

is upper triangular. Thus $\gamma c^2 - dcs - \beta s^2 = 0$, $d = \delta - \alpha$. Let $t = (d/2)^2 + \beta\gamma$; then $\cot \theta = (d/2 + \text{sign}(d) \sqrt{t})/2\gamma$, $s = \text{sign}(\cot \theta)(1 + \cot^2 \theta)^{-1/2}$, $c = s \cot \theta$.

(ii) It is also convenient to perform a supplementary plane rotation after a pair of complex conjugate eigenvalues, $\lambda \pm i\mu$, has been recorded in the course of the QR algorithm. In this case the transformation of the diagonal block is

$$\begin{pmatrix} c & s \\ -s & c \end{pmatrix} \begin{pmatrix} \alpha & \beta \\ \gamma & \delta \end{pmatrix} \begin{pmatrix} c & -s \\ s & c \end{pmatrix} = \begin{pmatrix} \lambda & \theta \\ \xi & \lambda \end{pmatrix}$$

where $\xi\theta = -\mu^2$. This device is not used in HQR2.

Note that it is not in general possible to transform

$$\begin{pmatrix} \alpha & \beta \\ \gamma & \delta \end{pmatrix} \rightarrow \begin{pmatrix} \lambda & -\mu \\ \mu & \lambda \end{pmatrix}$$

using *orthogonal* similarity transformations.

The purpose of the transformation is to yield a simple solution to certain systems of linear equations which must be solved. The supplementary plane rotation is done at the stage when the eigenvalues are being recorded in QR2NOZ. In this case $t \equiv p^2 + \beta\gamma < 0$, $p = (\alpha - \delta)/2$. We want to choose $c = \cos \theta$ and $s = \sin \theta$ so that

$$\alpha c^2 + (\beta + \gamma)cs + \delta s^2 = \delta c^2 - (\beta + \gamma)cs + \alpha s^2.$$

Hence

$$\tan 2\theta = 2sc/(c^2 - s^2) = -2p/\sigma = (2|p|/|\sigma|) \text{sign}(-p\sigma), \quad \sigma = \beta + \gamma.$$

Let $\tau = \sqrt{(\sigma^2 + 4p^2)}$. Then

$$\cos \theta = q = \sqrt{(\frac{1}{2}(1 + \cos 2\theta))} = \sqrt{((1 + |\sigma|/\tau)/2)},$$

$$\sin \theta = \sin 2\theta/2 \cos \theta = |p| \text{sign}(-p\sigma)/\tau q.$$

3.2 The Computation of the Eigenvectors of a Standardized Real, Block Upper Triangular Matrix. For each real eigenvalue λ the eigenvectors u, w^* satisfy $Tu = u\lambda$, $w^*T = \lambda w^*$.

For each complex conjugate pair of eigenvalues, $\lambda \pm i\mu$, the eigenvectors $u_1 \pm iu_2$, $w_1^* \mp iw_2^*$ satisfy $T(u_1, u_2) = (u_1, u_2)\hat{E}$, $(w_1, w_2)^*T = \hat{E}(w_1, w_2)^*$, where

$$\hat{E} = \begin{pmatrix} \lambda & \mu \\ -\mu & \lambda \end{pmatrix}.$$

In effecting the back substitution process in real arithmetic there are four different cases which can occur, depending on whether the matrices D and E shown below are 1×1 or 2×2 :

$$T = \begin{pmatrix} \cdot & \cdot & \cdot & \cdot \\ & D & \cdot & \cdot \\ & & \cdot & \cdot \\ & & O & E \\ & & & \cdot \end{pmatrix} \quad D = \begin{cases} \alpha \text{ or} \\ \begin{pmatrix} \alpha & \beta \\ \gamma & \alpha \end{pmatrix} \end{cases}$$

$$E = \begin{cases} \lambda \text{ or} \\ \begin{pmatrix} \lambda & \phi \\ \theta & \lambda \end{pmatrix} \end{cases}$$

where $\theta\phi = -\mu^2$. The positions of D and E should be exchanged when considering the row eigenvectors.

Type 1. pair-pair (E is 2×2 , D is 2×2). Imagine that the elements of u_1, u_2 in the same row as D are about to be computed. All elements below these have already been found, the elements below E being zero.

Let $j1, j2$ be the rows of T in which D lies. Then the unknowns are

$$V = \begin{pmatrix} u_1(j1) & u_2(j1) \\ u_1(j2) & u_2(j2) \end{pmatrix}.$$

The equation to be solved in the column case is $-DV + V\hat{E} = R$, where

$$R = \begin{pmatrix} r_1(j1) & r_2(j1) \\ r_1(j2) & r_2(j2) \end{pmatrix}, \quad r_\nu(m) = \sum_{k=j2+1}^{i+1} t_{m,k}u_\nu(k), \quad \begin{cases} m = j1, j2, \\ \nu = 1, 2. \end{cases} \quad (10)$$

In the row case let

$$V = \begin{pmatrix} w_1(j1) & w_2(j1) \\ w_1(j2) & w_2(j2) \end{pmatrix};$$

then the equations to be solved are

$$-V^T D + \hat{E} V^T = R^T, \quad (11)$$

where R is as above except that k runs from i to $j1 - 1$. Transposing yields

$$-D^T V + V \hat{E}^T = R.$$

Comparing this with the column case we see that it is only necessary to transpose D and \hat{E} (i.e. to exchange β, γ and $\mu, -\mu$) in order to use the same code for both cases.

The way in which these four linear equations in four unknowns are solved is described in Section 3.3.

Type 2. pair-single (E is 2×2 , D is 1×1). The relevant equations are

$$-\alpha(u_1(j), u_2(j)) + (u_1(j), u_2(j))\hat{E} = (r_1(j), r_2(j))$$

and

$$-\alpha(w_1(j), w_2(j)) + (w_1(j), w_2(j))\hat{E}^T = (r_1(j), r_2(j)).$$

Let $d = \lambda - \alpha$, $\text{den} = d^2 + \mu^2$, $\text{val} = \begin{cases} \mu & (\text{for column}) \\ -\mu & (\text{for row}) \end{cases}$. The solution for both cases is

$$\begin{aligned} v_1 &= (r_1 \cdot d + r_2 \cdot \text{val})/\text{den}, \\ v_2 &= (-r_1 \cdot \text{val} + r_2 \cdot d)/\text{den}. \end{aligned} \quad (12)$$

Type 3. single-pair (E is 1×1 , D is 2×2). The relevant equations are

$$-D \begin{pmatrix} u_1(j1) \\ u_1(j2) \end{pmatrix} + \begin{pmatrix} u_1(j1) \\ u_1(j2) \end{pmatrix} \lambda = \begin{pmatrix} r_1(j1) \\ r_1(j2) \end{pmatrix}$$

and the same equation for w_1 with D^T in place of D . Set $d = \lambda - \alpha$, $\text{den} = d^2 - \beta\gamma$. The solution is

$$\begin{aligned} v_1 &= (r_1(j1) \cdot d + r_1(j2) \cdot \tilde{\beta})/\text{den}, \\ v_2 &= (r_1(j1) \cdot \tilde{\gamma} + r_1(j2) \cdot d)/\text{den}, \end{aligned} \quad (13)$$

where $\tilde{\beta} = \begin{cases} \beta & (\text{for column}) \\ \gamma & (\text{for row}) \end{cases}$ and $\tilde{\gamma} = \begin{cases} \gamma & (\text{for column}) \\ \beta & (\text{for row}) \end{cases}$. In practice $\tilde{\beta} = T(JJ, J)$, $\tilde{\gamma} = T(J, JJ)$ and the indices J and JJ are given the values of $J1$ and $J2$ in the appropriate order.

Type 4. single-single (E is 1×1 , D is 1×1). $v_1 = r_1(j)/\text{den}$, $\text{den} = \lambda - \alpha$.

Type 5. formula breakdown. If in any of the previous cases $D = E$ then the formulas for solution break down. There are two cases to consider.

(i) Linear Independence. Any element v_j for which the formula yields $0/0$ can be set to any value, the most convenient is 0 . This represents the existence of a whole space of eigenvectors associated with E .

(ii) Defective Case. Any element v_j for which the formula yields a value exceeding $1/\text{TOL}$ will cause the condition number to exceed $1/\text{TOL}$. If this case is detected computation is interrupted, the condition number is set to $1/\text{TOL}$, and the program proceeds to the next eigenvalue.

We propose that $\text{TOL} = 10^{-30}$ will be suitable for most applications and most computers.

These tests make the code simple and machine independent. However, it is possible to devise matrices for which the given value $1/\text{TOL}$ for the condition is very misleading. We know of no failsafe procedure which does not involve deciding the rank of $T - \xi$ for all ξ in a neighborhood of λ . This is a costly, difficult, and often unrewarding task.

3.3 *Closed Form Solution for Equations of Type 1.* The equations to be solved are of the form $-DV + V\hat{E} = R$, where

$$D = \begin{pmatrix} \alpha & \beta \\ \gamma & \alpha \end{pmatrix}, \quad \hat{E} = \begin{pmatrix} \lambda & \mu \\ -\mu & \lambda \end{pmatrix}, \quad R = \begin{pmatrix} r_{11} & r_{12} \\ r_{21} & r_{22} \end{pmatrix}.$$

The standardization of the block triangular matrix T forces the diagonal elements of D to be equal. This yields simple formulas for the elements of V .

Rewrite the equation as

$$(v_{11} \ v_{12} \ v_{21} \ v_{22}) \begin{pmatrix} B & -\gamma I_2 \\ -\beta I_2 & B \end{pmatrix} = (r_{11} \ r_{12} \ r_{21} \ r_{22}) = r^T,$$

where

$$B = \begin{pmatrix} \lambda - \alpha & \mu \\ -\mu & \lambda - \alpha \end{pmatrix}.$$

Observe that

$$\begin{pmatrix} B & -\gamma I_2 \\ -\beta I_2 & B \end{pmatrix} \begin{pmatrix} B^T & \gamma I_2 \\ \beta I_2 & B^T \end{pmatrix} = \tau I_4 + 2\mu J_4,$$

where $\tau = (\lambda - \alpha)^2 + \mu^2 - \beta\gamma$, with $\beta\gamma < 0$,

$$J_4 = \begin{pmatrix} 0 & 0 & 0 & \gamma \\ 0 & 0 & -\gamma & 0 \\ 0 & \beta & 0 & 0 \\ -\beta & 0 & 0 & 0 \end{pmatrix},$$

and I_k is the $k \times k$ identity matrix. Further, $(\tau I_4 + 2\mu J_4)(\tau I_4 - 2\mu J_4) = [\tau^2 - 4\mu^2(-\beta\gamma)]I_4$. Hence

$$\begin{aligned} (v_{11} \ v_{12} \ v_{21} \ v_{22})(\tau^2 + 4\mu^2\beta\gamma) &= r^T \begin{pmatrix} B^T & \gamma I_2 \\ \beta I_2 & B^T \end{pmatrix} (\tau I_4 - 2\mu J_4) \\ &= r^T \begin{pmatrix} e & -f & \gamma g & -\gamma h \\ f & e & \gamma h & \gamma g \\ \beta g & -\beta h & e & -f \\ \beta h & \beta g & f & e \end{pmatrix} \end{aligned}$$

where $d = \lambda - \alpha$, $e = d\tau$, $f = \mu(\tau + 2\beta\gamma)$, $g = \tau - 2\mu^2$, and $h = 2d\mu$. Note that

$$\begin{aligned} \tau^2 + 4\mu^2\beta\gamma &= (d^2 + \mu^2 - \beta\gamma)^2 + 4\mu^2\beta\gamma \\ &= g^2 + h^2 \\ &= 0 \quad \text{if and only if} \quad \alpha = \lambda, \mu^2 = -\beta\gamma. \end{aligned}$$

These same formulas will be valid for the row eigenvectors provided we exchange (β, γ) and $(\mu, -\mu)$.

The alternative to using this closed form solution is to code a special version of Gaussian elimination with pivoting. It is the pivoting which would lengthen the code considerably.

3.4 The Condition Number of Conjugate Pairs of Eigenvalues. Let $\lambda \pm i\mu$ be a complex pair of eigenvalues of the real Schur matrix T obtained by the QR algorithm. In the course of the algorithm the following real equations are solved for real n -vectors u_1, u_2, w_1, w_2 :

$$T(u_1, u_2) = (u_1, u_2) \begin{pmatrix} \rho & \mu \\ -\mu & \rho \end{pmatrix}, \quad (w_1, w_2)^* T = \begin{pmatrix} \rho & \mu \\ -\mu & \rho \end{pmatrix} (w_1, w_2)^*. \quad (5)$$

Thus $\text{span}(u_1, u_2)$ and $\text{span}(w_1^*, w_2^*)$ are real invariant subspaces under T . However, $\{u_1, u_2\}$ and $\{w_1^*, w_2^*\}$ are very special bases of these spaces.

LEMMA. *With the notation given above, $u_1 \pm iu_2$ and $w_1^* \mp iw_2^*$ are the column and row eigenvectors belonging to $\lambda \pm i\mu$.*

PROOF. From (5),

$$\begin{aligned} Tu_1 &= u_1\lambda - u_2\mu, & w_1^* T &= \lambda w_1^* + \mu w_2^*, \\ Tu_2 &= u_1\mu + u_2\lambda, & w_2^* T &= -\mu w_1^* + \lambda w_2^*. \end{aligned}$$

Hence

$$T(u_1 + iw_2) = u_1(\lambda + i\mu) + iw_2(i\mu + \lambda),$$

then the equations to be solved are

$$-V^T D + \hat{E} V^T = R^T, \quad (11)$$

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$$\begin{aligned} v_1 &= (r_1 \cdot d + r_2 \cdot \text{val})/\text{den}, \\ v_2 &= (-r_1 \cdot \text{val} + r_2 \cdot d)/\text{den}. \end{aligned} \quad (12)$$

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and the same equation for w_1 with D^T in place of D . Set $d = \lambda - \alpha$, $\text{den} = d^2 - \beta\gamma$. The solution is

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where $\tilde{\beta} = \begin{cases} \beta & \text{(for column)} \\ \gamma & \text{(for row)} \end{cases}$ and $\tilde{\gamma} = \begin{cases} \gamma & \text{(for column)} \\ \beta & \text{(for row)} \end{cases}$. In practice $\tilde{\beta} = T(JJ, J)$, $\tilde{\gamma} = T(J, JJ)$ and the indices J and JJ are given the values of $J1$ and $J2$ in the appropriate order.

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$$(v_{11} \ v_{12} \ v_{21} \ v_{22}) \begin{pmatrix} B & -\gamma I_2 \\ -\beta I_2 & B \end{pmatrix} = (r_{11} \ r_{12} \ r_{21} \ r_{22}) = r^T,$$

where

$$B = \begin{pmatrix} \lambda - \alpha & \mu \\ -\mu & \lambda - \alpha \end{pmatrix}.$$

Observe that

$$\begin{pmatrix} B & -\gamma I_2 \\ -\beta I_2 & B \end{pmatrix} \begin{pmatrix} B^T & \gamma I_2 \\ \beta I_2 & B^T \end{pmatrix} = \tau I_4 + 2\mu J_4,$$

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$$T(u_1, u_2) = (u_1, u_2) \begin{pmatrix} \rho & \mu \\ -\mu & \rho \end{pmatrix}, \quad (w_1, w_2)^* T = \begin{pmatrix} \rho & \mu \\ -\mu & \rho \end{pmatrix} (w_1, w_2)^*. \quad (5)$$

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$$\begin{aligned} Tu_1 &= u_1\lambda - u_2\mu, & w_1^* T &= \lambda w_1^* + \mu w_2^*, \\ Tu_2 &= u_1\mu + u_2\lambda, & w_2^* T &= -\mu w_1^* + \lambda w_2^*. \end{aligned}$$

Hence

$$T(u_1 + iu_2) = u_1(\lambda + i\mu) + iu_2(i\mu + \lambda),$$

$$(w_1^* - iw_2^*)T = (\lambda + i\mu)w_1^* - (\lambda + i\mu)iw_2^*.$$

The eigenvectors for $\lambda - i\mu$ are obtained in the same way. \square

Consequently,

$$\text{cond}(\lambda \pm i\mu) = \|u_1 + iu_2\| \cdot \|w_1^* - iw_2^*\| / [w_1^*u_1 + w_2^*u_2 + i(w_1^*u_2 - w_2^*u_1)].$$

Use was made in the lemma of the quasi-triangular nature of T . A consequence of this form is that u_1 and w_1 can be packed into the same real n -vector with two overlapping elements as indicated:

$$\left. \begin{aligned} u_1^* &= (x, \dots, x, p_i, q_i, 0, \dots, 0) \\ w_i^* &= (0, \dots, 0, \bar{p}_i, \bar{q}_i, x, \dots, x) \end{aligned} \right\} \quad i = 1, 2.$$

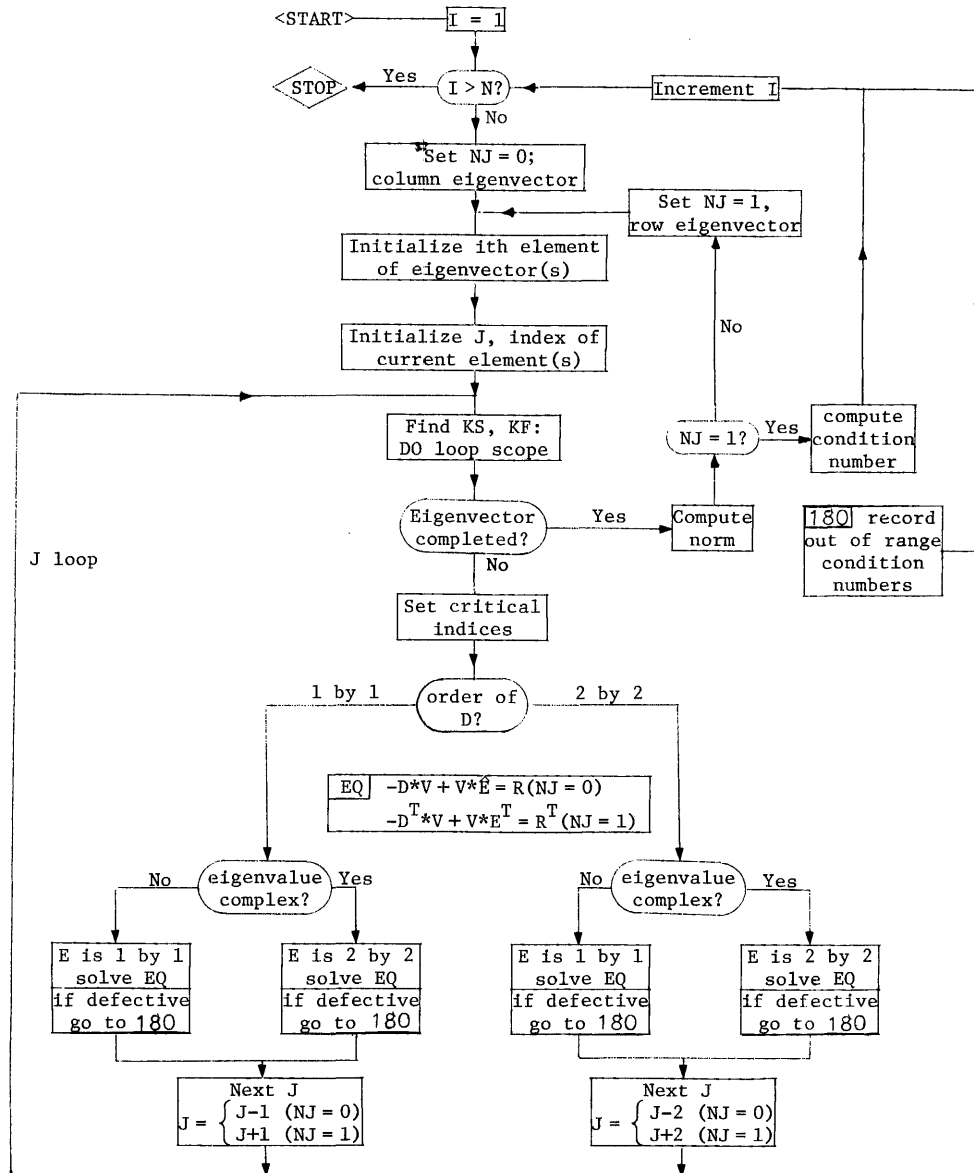


Fig. 1. Flowchart representing way in which column and row eigenvectors are computed allowing for haphazard ordering of 1x1 and 2x2 blocks on diagonal of real Schur form

The equations to be satisfied by p_i, q_i are of the form

$$\begin{pmatrix} \lambda & \beta \\ \gamma & \lambda \end{pmatrix} \begin{pmatrix} p_1 & p_2 \\ q_1 & q_2 \end{pmatrix} - \begin{pmatrix} p_1 & p_2 \\ q_1 & q_2 \end{pmatrix} \begin{pmatrix} \lambda & \mu \\ -\mu & \lambda \end{pmatrix} = 0,$$

$$\begin{pmatrix} \bar{p}_1 & \bar{q}_1 \\ \bar{p}_2 & \bar{q}_2 \end{pmatrix} \begin{pmatrix} \lambda & \beta \\ \gamma & \lambda \end{pmatrix} - \begin{pmatrix} \lambda & \mu \\ -\mu & \lambda \end{pmatrix} \begin{pmatrix} \bar{p}_1 & \bar{q}_1 \\ \bar{p}_2 & \bar{q}_2 \end{pmatrix} = 0,$$

where $\mu^2 = -\beta\gamma$. These equations reduce to

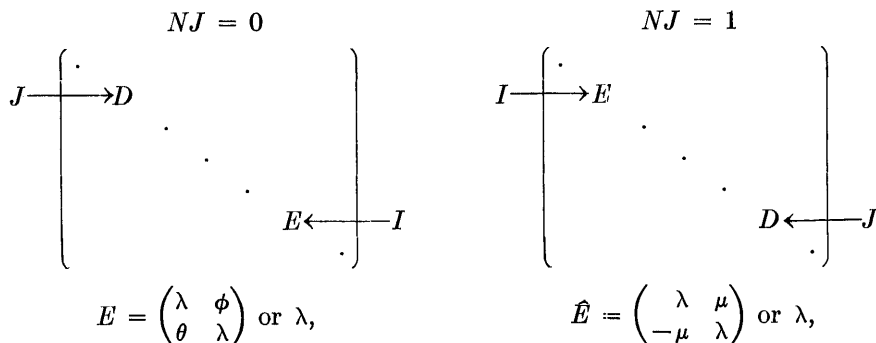
$$\begin{aligned} \beta q_2 &= p_1 \mu, & \bar{p}_1 \beta &= \mu \bar{q}_2, \\ \beta q_1 &= -\mu p_2, & \beta \bar{p}_2 &= -\mu \bar{q}_1. \end{aligned}$$

The simplest solution (which we adopt) takes $p_1 = \bar{p}_1 = 1, q_1 = p_2 = \bar{q}_1 = \bar{p}_2 = 0,$
 $q_2 = \mu/\beta, \bar{q}_2 = 1/q_2$. With this choice,

$$\begin{aligned} (w_1 - iw_2)^*(u_1 + iu_2) &= (\bar{p}_1 p_1 + \bar{p}_2 p_2) + (\bar{q}_1 q_1 + \bar{q}_2 q_2) = 2 \text{ and} \\ \text{cond}(\lambda \pm i\mu) &= [(\|u_1\|^2 + \|u_2\|^2)(\|w_1\|^2 + \|w_2\|^2)]^{1/2}/2. \end{aligned} \tag{14}$$

4. The Structure of CONDIR

As indicated in Section 3.2 the same section of code (the J loop) computes the column and the row eigenvector of the I th eigenvalue. It is the standardization of the Schur form which permits this economy. J , which always points to the block D , decreases for the column eigenvector ($NJ = 0$) and increases for the row eigenvector ($NJ = 1$), as indicated:



where $\theta\phi = -\mu^2$.

Is the 0 diagonal element in I^0 a keypunch error or did it really belong in the user's problem?

$$\begin{aligned} L^0 &= \begin{pmatrix} 0 & X \\ I^0 & 0 \end{pmatrix}; \quad I^0 = \text{diag}(0, 1, 1, \dots, 1); \\ X &= \begin{pmatrix} D & T_1 & 0 \\ T_2 & F_1 & -F_2 \\ 0 & -F_3 & F_4 \end{pmatrix} 10^8; \quad F = \begin{pmatrix} a & 0 & 0 & b \\ 0 & a & -b & 0 \\ 0 & -c & d & 0 \\ c & 0 & 0 & d \end{pmatrix}; \\ D &= \text{diag}(.5221, .3563, .5552 \times 10^{-3}, .1328); \\ T_1 &= \begin{pmatrix} .5221 & 0 & 0 & -.8951 \\ 0 & -.3563 & .6109 & 0 \\ 0 & 0 & -.5552 \times 10^{-3} & 0 \\ 0 & 0 & 0 & -.1328 \end{pmatrix}; \\ T_2 &= \begin{pmatrix} -.0976 & 0 & 0 & 0 \\ 0 & -.0666 & 0 & 0 \\ 0 & .02659 & -.4218 \times 10^{-4} & 0 \\ -.03896 & 0 & 0 & -.1009 \times 10^{-1} \end{pmatrix}; \\ F_1 &= \begin{pmatrix} 2.859 & 0 & 0 & 1.079 \\ 0 & 2.828 & -1.026 & 0 \\ 0 & -.2389 & .4294 & 0 \\ .2513 & 0 & 0 & .4607 \end{pmatrix}; \\ F_2 &= F \begin{pmatrix} 2.761 & .5891 \\ .2123 & .3590 \end{pmatrix}; \quad F_3 = F \begin{pmatrix} 1.627 & .5373 \\ .1096 & .2868 \end{pmatrix}; \quad F_4 = F \begin{pmatrix} 1.849 & .3731 \\ .1178 & .4970 \end{pmatrix}; \\ \|L^0\|_\infty &\approx 5 \times 10^8. \end{aligned}$$

Fig. 2. 24x24 matrix for case study

Table I. Eigenvalues and Condition Numbers of \mathbf{L}^0 , \mathbf{L} , \mathbf{M}

$\lambda_i(\mathbf{L}^0)$	Cond(\mathbf{L})	Cond(\mathbf{L}^0) (balanced)	$\lambda_i(\mathbf{L})$	Cond(\mathbf{L})	$\lambda_i(\mathbf{M})$	Cond(\mathbf{M})
$\pm .21534594 \times 10^5$	10^4	1	$\pm .21534594 \times 10^5$	10^4	$\pm .21534594 \times 10^5$	10^4
$\pm .18667890 \times 10^5$	10^4	10^3	$\pm .18654343 \times 10^5$	10^4	$\pm .18692513 \times 10^5$	10^4
$\pm .11076317 \times 10^{5i}$	6×10^3	4×10^3	$\pm .1098663V \times 10^{5i}$	6×10^3	$\pm .11142610 \times 10^{5i}$	6×10^3
$\pm .82614552 \times 10^4$	6×10^3	5×10^3	$\pm .8646568V \times 10^4$	10^4	$\pm .86339960 \times 10^{4*}$	10^4
$\pm .83281998 \times 10^4$	6×10^3	1	$\pm .83281998 \times 10^4$	6×10^3	$\pm .83281998 \times 10^4$	6×10^3
$\pm .41438248 \times 10^4$	7×10^3	10^4	$\pm .7026706V \times 10^4$	10^4	$\pm .71883679 \times 10^{4*}$	10^4
$\pm .64209960 \times 10^4$	7×10^3	3	$\pm .64209960 \times 10^4$	6×10^3	$\pm .64209960 \times 10^4$	7×10^3
$\pm .33632953 \times 10^4$	5×10^3	9×10^3	$\pm .38448590 \times 10^4$	3×10^3	$\pm .38458962 \times 10^{4*}$	3×10^3
$\pm .35102700 \times 10^4$	2×10^3	4	$\pm .35102700 \times 10^4$	2×10^3	$\pm .35102700 \times 10^4$	2×10^3
$\pm .30530592 \times 10^4$	3×10^3	3	$\pm .30530592 \times 10^4$	3×10^3	$\pm .30530592 \times 10^4$	3×10^3
$\pm .0\dagger$	10^{15}	10^{30}	$\pm .29241979 \times 10^4$	3×10^3	$\pm .29134631 \times 10^4$	3×10^3
$\pm .23559292 \times 10^3$	10^2	1	$\pm .23559291 \times 10^3$	10^2	$\pm .23559292 \times 10^3$	10^2

* The imaginary pair of eigenvalues had real parts less than 10^{-6} (a relative error of 10^{-11}). V denotes a digit that changed when the matrix was balanced.

† The unbalanced matrix \mathbf{L}^0 had a negative eigenvalue -2×10^{-8} instead of -0 .

At each major step in the J loop the equations EQ, given in the chart, are solved as D varies. Four cases are shown in the flowchart which is Figure 1.

5. Results

The matrix \mathbf{L}^0 described in Figure 2 came (in punched card form) from a large industrial company. It was causing their eigenvalue program to fail.

An inspection of the form of \mathbf{L}^0 suggests that the strange diagonal element in \mathbf{I}^0 and the discordant sign of the (1, 1) element of \mathbf{T}_1 were keypunch errors. Let us consider the matrices resulting from the removal of these anomalies:

$$\mathbf{L} = \begin{pmatrix} \mathbf{0} & \mathbf{X} \\ \mathbf{I} & \mathbf{0} \end{pmatrix}, \quad \mathbf{M} = \begin{pmatrix} \mathbf{0} & \mathbf{Y} \\ \mathbf{I} & \mathbf{0} \end{pmatrix},$$

$$\mathbf{Y} = 10^8 \begin{pmatrix} \mathbf{D} & \mathbf{T}_1^- & \mathbf{0} \\ \mathbf{T}_2 & \mathbf{T}_1 & -\mathbf{F}_2 \\ \mathbf{0} & -\mathbf{F}_3 & \mathbf{F}_4 \end{pmatrix},$$

where \mathbf{T}_1^- is obtained from \mathbf{T}_1 by reversing the sign of its (1, 1) element.

Notice that the eigenvalues of \mathbf{L} are the square roots of those of \mathbf{X} :

$$\begin{pmatrix} \mathbf{0} & \mathbf{X} \\ \mathbf{I} & \mathbf{0} \end{pmatrix} \begin{pmatrix} \mathbf{u} \\ \mathbf{v} \end{pmatrix} = \lambda \begin{pmatrix} \mathbf{u} \\ \mathbf{v} \end{pmatrix} \leftrightarrow \mathbf{X}\mathbf{v} = \lambda^2\mathbf{v}, \quad \mathbf{u} = \lambda\mathbf{v}.$$

The eigenvalues of \mathbf{L}^0 , \mathbf{L} , and \mathbf{M} are given in Table I and we offer the following comments. Every eigenvalue of \mathbf{L}^0 is moderately ill conditioned and the zero pair appear to belong to a quadratic elementary divisor (only one eigenvector). Perhaps some of this is due to the unbalanced nature of \mathbf{L}^0 . The thirteenth row of \mathbf{L}^0 is null and this must be permuted out of the way before the rest of the matrix can

be balanced. The result is that none of the computed eigenvalues changed, but half of them became almost perfectly conditioned.

In fact we can say that the ill condition of all six pairs is due to the zero element in position (13, 1). When this is replaced by one, we obtain the matrix \mathbf{L} which has six pairs of eigenvalues almost identical to the well conditioned pairs of the balanced \mathbf{L}^0 . Four of the other six pairs are changed completely, the remaining two ($\pm .186 \times 10^6$ and $\pm .11 \times 10^6 i$) are substantially altered. Interestingly the balanced versions of \mathbf{L} and \mathbf{M} are almost normal and we have not bothered to record the condition numbers. The six pairs of eigenvalues which were unchanged by the move from \mathbf{L}^0 to \mathbf{L} were also invariant in the change from \mathbf{L} to \mathbf{M} . The other six pairs had relative errors less than 2.5 percent.

We can tell in advance what the balanced form of \mathbf{L} and \mathbf{M} will be:

$$\hat{\mathbf{L}} = \begin{pmatrix} 0 & 10^{-4}X \\ 10^4I & 0 \end{pmatrix}, \quad \hat{\mathbf{M}} = \begin{pmatrix} 0 & 10^{-4}Y \\ 10^4I & 0 \end{pmatrix}.$$

The change from \mathbf{L}^0 to \mathbf{L} is tiny relative to $\|\mathbf{L}^0\|$ ($\doteq 10^{-8} \|\mathbf{L}^0\|$) but the change from $\hat{\mathbf{L}}^0$ to $\hat{\mathbf{L}}$ is approximately $\|\hat{\mathbf{L}}\|$.

We conclude that the suspicious element in \mathbf{L}^0 was probably a keypunch error. Concerning the (1, 1) element of \mathbf{T}_1 we cannot say, both \mathbf{L} and \mathbf{M} are reasonable matrices and indeed the change of sign does not affect the leading two decimals in any eigenvalue.

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ALGORITHM

SUBROUTINE QR2NOZ(NM, N, LOW, IGH, H, WR, WI, IERR)	QR2	10
C PURPOSE	QR2	20
C THE FORTRAN SUBROUTINE QR2NOZ COMPUTES THE EIGENVALUES OF A REAL	QR2	30
C UPPER HESSENBERG MATRIX USING THE QR METHOD AND REDUCES THE	QR2	40
C MATRIX TO A STANDARDIZED QUASI-TRIANGULAR FORM. COMPUTATIONS	QR2	50
C ARE DONE IN REAL ARITHMETIC.	QR2	60
C THE SUBROUTINE STATEMENT IS	QR2	70
C SUBROUTINE QR2NOZ(NM,N,LOW,IGH,H,WR,WI,IERR).	QR2	80
C NM IS AN INTEGER INPUT VARIABLE SET EQUAL TO THE ROW DIMENSION	QR2	90
C OF THE ARRAY H AS SPECIFIED IN THE CALLING PROGRAM.	QR2	100
C N IS AN INTEGER INPUT VARIABLE SET EQUAL TO THE ORDER OF THE	QR2	110
C MATRIX H. N .LE. NM	QR2	120
C LOW,IGH ARE INTEGER INPUT VARIABLES INDICATING THE BOUNDARY INDICES	QR2	130
C FOR THE BALANCED MATRIX. IF THE MATRIX IS NOT BALANCED SET	QR2	140
C LOW TO 1 AND IGH TO N.	QR2	150
C H IS A REAL TWO-DIMENSIONAL ARRAY WITH ROW DIMENSION NM AND	QR2	160
C COLUMN DIMENSION AT LEAST N. ON INPUT IT CONTAINS THE	QR2	170
C UPPER HESSENBERG MATRIX OF ORDER N. ON OUTPUT IT CONTAINS	QR2	180
C THE STANDARDIZED QUASI-TRIANGULAR MATRIX.	QR2	190
C WR,WI ARE REAL OUTPUT ONE-DIMENSIONAL VARIABLES OF DIMENSION AT	QR2	200
C LEAST N CONTAINING THE REAL AND IMAGINARY PARTS,	QR2	210
C RESPECTIVELY, OF THE EIGENVALUES OF THE HESSENBERG MATRIX.	QR2	220
C THE EIGENVALUES ARE UNORDERED EXCEPT THAT COMPLEX CONJUGATE	QR2	230
C PAIRS OF EIGENVALUES APPEAR CONSECUTIVELY WITH THE	QR2	240
C EIGENVALUE HAVING THE POSITIVE IMAGINARY PART FIRST.	QR2	250


```

100 MP2 = M + 2
DO 110 I=MP2,EN
  H(I,I-2) = 0.0
  IF (I.EQ.MP2) GO TO 110
  H(I,I-3) = 0.0
110 CONTINUE
C DOUBLE QR STEP INVOLVING ROWS L TO EN
C AND COLUMNS M TO EN.
DO 190 K=M,NA
  NOTLAS = K.NE.NA
  IF (K.EQ.M) GO TO 120
  P = H(K,K-1)
  Q = H(K+1,K-1)
  R = 0.0
  IF (NOTLAS) R = H(K+2,K-1)
  X = ABS(P) + ABS(Q) + ABS(R)
  IF (X.EQ.0.0) GO TO 190
  P = P/X
  Q = Q/X
  R = R/X
120 S = SIGN(SQRT(P*P+Q*Q+R*R),P)
  IF (K.EQ.M) GO TO 130
  H(K,K-1) = -S*X
  GO TO 140
130 IF (L.NE.M) H(K,K-1) = -H(K,K-1)
140 P = P + S
  X = P/S
  Y = Q/S
  ZZ = R/S
  Q = Q/P
  R = R/P
C ROW MODIFICATION
DO 160 J=K,N
  P = H(K,J) + Q*H(K+1,J)
  IF (.NOT.NOTLAS) GO TO 150
  P = P + R*H(K+2,J)
  H(K+2,J) = H(K+2,J) - P*ZZ
150 H(K+1,J) = H(K+1,J) - P*Y
  H(K,J) = H(K,J) - P*X
160 CONTINUE
  J = MIN0(EN,K+3)
C COLUMN MODIFICATION
DO 180 I=1,J
  P = X*H(I,K) + Y*H(I,K+1)
  IF (.NOT.NOTLAS) GO TO 170
  P = P + ZZ*H(I,K+2)
  H(I,K+2) = H(I,K+2) - P*R
170 H(I,K+1) = H(I,K+1) - P*Q
  H(I,K) = H(I,K) - P
180 CONTINUE
190 CONTINUE
GO TO 30
C ONE ROOT FOUND
200 H(EN,EN) = X + T
  WR(EN) = H(EN,EN)
  WI(EN) = 0.0
  EN = NA
  GO TO 20
C TWO ROOTS FOUND
210 P = (Y-X)/2.0
  Q = P*P + W
  ZZ = SQRT(ABS(Q))
  H(EN,EN) = X + T
  X = H(EN,EN)
  H(NA,NA) = Y + T
  IF (Q.LT.0.0) GO TO 220
  ZZ = P + SIGN(ZZ,P)
C REAL PAIR
  WR(NA) = X + ZZ
  WR(EN) = WR(NA)
  IF (ZZ.NE.0.0) WR(EN) = X - W/ZZ
  WI(NA) = 0.0
  WI(EN) = 0.0
  X = H(EN,NA)
  R = SQRT(X*X+ZZ*ZZ)
  P = X/R
QR2 1020
QR2 1030
QR2 1040
QR2 1050
QR2 1060
QR2 1070
QR2 1080
QR2 1090
QR2 1100
QR2 1110
QR2 1120
QR2 1130
QR2 1140
QR2 1150
QR2 1160
QR2 1170
QR2 1180
QR2 1190
QR2 1200
QR2 1210
QR2 1220
QR2 1230
QR2 1240
QR2 1250
QR2 1260
QR2 1270
QR2 1280
QR2 1290
QR2 1300
QR2 1310
QR2 1320
QR2 1330
QR2 1340
QR2 1350
QR2 1360
QR2 1370
QR2 1380
QR2 1390
QR2 1400
QR2 1410
QR2 1420
QR2 1430
QR2 1440
QR2 1450
QR2 1460
QR2 1470
QR2 1480
QR2 1490
QR2 1500
QR2 1510
QR2 1520
QR2 1530
QR2 1540
QR2 1550
QR2 1560
QR2 1570
QR2 1580
QR2 1590
QR2 1600
QR2 1610
QR2 1620
QR2 1630
QR2 1640
QR2 1650
QR2 1660
QR2 1670
QR2 1680
QR2 1690
QR2 1700
QR2 1710
QR2 1720
QR2 1730
QR2 1740
QR2 1750
QR2 1760
QR2 1770

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      Q = ZZ/R                                QR2 1780
      GO TO 230                                QR2 1790
C COMPLEX PAIR                                QR2 1800
      220 WR(NA) = X + P                       QR2 1810
      WR(EN) = X + P                           QR2 1820
      WI(NA) = ZZ                               QR2 1830
      WI(EN) = -ZZ                              QR2 1840
C MAKE DIAGONAL ELEMENTS EQUAL              QR2 1850
      IF (P.EQ.0.0) GO TO 260                  QR2 1860
      BPC = H(EN,NA) + H(NA,EN)                QR2 1870
      TX = SQRT(BPC*BPC+4.0*P*P)               QR2 1880
      Q = SQRT(.5*(1.0+ABS(BPC)/TX))           QR2 1890
      P = SIGN(P/(Q*TX),-BPC*P)               QR2 1900
C ROW MODIFICATION                           QR2 1910
      230 DO 240 J=NA,N                       QR2 1920
          ZZ = H(NA,J)                         QR2 1930
          H(NA,J) = Q*ZZ + P*H(EN,J)           QR2 1940
          H(EN,J) = Q*H(EN,J) - P*ZZ          QR2 1950
      240 CONTINUE                            QR2 1960
C COLUMN MODIFICATION                        QR2 1970
      DO 250 I=1,N                             QR2 1980
          ZZ = H(I,NA)                         QR2 1990
          H(I,NA) = Q*ZZ + P*H(I,EN)           QR2 2000
          H(I,EN) = Q*H(I,EN) - P*ZZ          QR2 2010
      250 CONTINUE                            QR2 2020
      260 EN = ENM2                             QR2 2030
      GO TO 200                                QR2 2040
      270 IERR = EN                             QR2 2050
      RETURN                                    QR2 2060
      END                                       QR2 2070

      SUBROUTINE CONDIT(NM, N, A, V1, V2, WI, COND)    CON 10
C CONDIT COMPUTES THE CONDITION NUMBERS OF THE EIGENVALUES OF A    CON 20
C STANDARDIZED QUASI-TRIANGULAR MATRIX.                CON 30
C THE SUBROUTINE STATEMENT IS                            CON 40
C SUBROUTINE CONDIT(NM,N,A,V1,V2,WI,COND).              CON 50
C ON INPUT                                                CON 60
C NM MUST BE SET TO THE ROW DIMENSION OF THE TWO DIMENSIONAL    CON 70
C ARRAY AS DECLARED IN THE CALLING PROGRAM.              CON 80
C N IS THE ORDER OF THE MATRIX. N.LE.NM                  CON 90
C A CONTAINS THE STANDARDIZED QUASI-TRIANGULAR MATRIX PRODUCED    CON 100
C BY QR2NOZ.                                             CON 110
C WI CONTAINS THE IMAGINARY PARTS OF THE EIGENVALUES. THE    CON 120
C EIGENVALUES ARE UNORDERED EXCEPT THAT COMPLEX CONJUGATE    CON 130
C PAIRS APPEAR CONSECUTIVELY.                            CON 140
C V1,V2 ARE FOR TEMPORARY STORAGE.                       CON 150
C ON OUTPUT                                              CON 160
C A IS UNALTERED.                                       CON 170
C COND CONTAINS THE CONDITION NUMBERS CORRESPONDING TO THE    CON 180
C EIGENVALUES IN (V2,WI). COND = 1./TOL IF THE USUAL    CON 190
C FORMULA WOULD CAUSE OVERFLOW OR YIELD A VALUE EXCEEDING    CON 200
C 1/TOL. TOL NEED NOT DEPEND ON THE COMPUTER.            CON 210
C V2 CONTAINS THE REAL PARTS OF THE EIGENVALUES.         CON 220
C TYPICAL USAGE                                         CON 230
C DIMENSION A(50,50),WR(50),WI(50),COND(50),ORT(50)      CON 240
C *****ENTER MATRIX A AND DIMENSIONS N,NM*****CON 250
C LOW=1                                                  CON 260
C IGH=N                                                  CON 270
C CALL ORTHES(NM,N,LOW,IGH,A,ORT)                       CON 280
C CALL QR2NOZ(NM,N,LOW,IGH,A,WR,WI,IERR)                CON 290
C CALL CONDIT(NM,N,A,ORT,WR,WI,COND)                    CON 300
C *****CON 310
C NOTE THE USE OF ORT AND WR IN CONDIT                  CON 320
C DIMENSION A(NM,NM), V1(NM), V2(NM), WI(NM), COND(NM)    CON 330
C DIMENSION R1(2), R2(2)                                CON 340
C DATA TOL /1.E-30/                                    CON 350
C -----CON 360
      I = 1                                           CON 370
      10 IF (I.GT.N) GO TO 190                        CON 380
      VALR = A(I,I)                                    CON 390
      VALI = WI(I)                                     CON 400
      VALI2 = VALI*VALI                                CON 410
C NJ GIVES EIGENVECTOR TYPE, 0 FOR COLUMN, 1 FOR ROW    CON 420
      NJ = 0                                           CON 430
C INITIALIZE NONZERO ELEMENTS OF EIGENVECTOR (V1,V2)    CON 440
      V1(I) = 1.0                                       CON 450

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V2(I) = 0.0 CON 460
20 J = I - 1 + 2*NJ CON 470
  IF (VALI.EQ.0.0) GO TO 30 CON 480
  V2(I+1) = VALI/A(I,I+1) CON 490
  V1(I+1) = 0.0 CON 500
  IF (NJ.EQ.1) V2(I+1) = 1.0/V2(I+1) CON 510
  J = I - 1 + 3*NJ CON 520
C FIND THE INDICES OF ELEMENTS COMPUTED SO FAR CON 530
30 KS = J + 1 + NJ*(I-J-1) CON 540
  KF = I + 1 + NJ*(J-I-2) CON 550
  IF (VALI.EQ.0.0 .AND. NJ.EQ.0) KF = KF - 1 CON 560
C TEST FOR COMPLETION OF EIGENVECTOR CON 570
  IF ((J+NJ.LT.1) .OR. (J+NJ.GT.N+1)) GO TO 130 CON 580
C ***** CON 590
C *SOLVE -D*V + V*E = R FOR V = (V1,V2). D IS A DIAGONAL BLOCK IN ROWS CON 600
C * J1,J2, AND E IS THE REAL CANONICAL FORM OF THE ITH EIGENVALUE. CON 610
C * EITHER D OR E OR BOTH CAN BE 1 BY 1 CON 620
C ***** CON 630
C FIND J1 AND J2 (J1.LE.J2) FOR ALL CASES CON 640
  JJ = J CON 650
  IF (WI(J).NE.0.0) JJ = J - 1 + 2*NJ CON 660
  J0 = NJ*(J-JJ) CON 670
  J1 = JJ + J0 CON 680
  J2 = J - J0 CON 690
  D1 = VALR - A(J,J) CON 700
C CALCULATE RIGHT HAND SIDE R CON 710
  DO 70 L=J1,J2 CON 720
    LJ = L - J1 + 1 CON 730
    R1(LJ) = 0.0 CON 740
    R2(LJ) = 0.0 CON 750
    IF (VALI.NE.0.0) GO TO 50 CON 760
    DO 40 K=KS,KF CON 770
      LK = NJ*(K-L) CON 780
      AA = A(L+LK,K-LK) CON 790
      R1(LJ) = R1(LJ) + AA*V1(K) CON 800
40 CONTINUE CON 810
  GO TO 70 CON 820
50 DO 60 K=KS,KF CON 830
  LK = NJ*(K-L) CON 840
  AA = A(L+LK,K-LK) CON 850
  R1(LJ) = R1(LJ) + AA*V1(K) CON 860
  R2(LJ) = R2(LJ) + AA*V2(K) CON 870
60 CONTINUE CON 880
70 CONTINUE CON 890
  IF (JJ.NE.J) GO TO 100 CON 900
C *****D IS 1 BY 1 ***** CON 910
  IF (VALI.NE.0.0) GO TO 80 CON 920
C E IS 1 BY 1 (D IS 1 BY 1) CON 930
  IF (ABS(D1).LT.TOL*ABS(R1(1))) GO TO 180 CON 940
  V1(J) = 0.0 CON 950
  V2(J) = 0.0 CON 960
  IF (D1.NE.0.0) V1(J) = R1(1)/D1 CON 970
  GO TO 90 CON 980
C E IS 2 BY 2 ( D IS 1 BY 1 ) CON 990
80 DEN = D1*D1 + VALI2 CON 1000
  VAL = VALI*(-1.0)**NJ CON 1010
  V1(J) = R1(1)*D1 + R2(1)*VAL CON 1020
  V2(J) = R2(1)*D1 - R1(1)*VAL CON 1030
  VMAX = AMAX1(ABS(V1(J)),ABS(V2(J))) CON 1040
  IF (DEN.LT.TOL*VMAX) GO TO 180 CON 1050
  V1(J) = V1(J)/DEN CON 1060
  V2(J) = V2(J)/DEN CON 1070
C NEXT J CON 1080
90 J = J - 1 + 2*NJ CON 1090
  GO TO 30 CON 1100
C *****D IS 2 BY 2***** CON 1110
100 IF (VALI.NE.0.0) GO TO 110 CON 1120
C E IS 1 BY 1 (D IS 2 BY 2) CON 1130
  DEN = D1*D1 + WI(J)**2 CON 1140
  V2(J1) = 0.0 CON 1150
  V2(J2) = 0.0 CON 1160
  V1(J1) = R1(1)*D1 + R1(2)*A(JJ,J) CON 1170
  V1(J2) = R1(1)*A(J,JJ) + R1(2)*D1 CON 1180
  VMAX = AMAX1(ABS(V1(J1)),ABS(V1(J2))) CON 1190
  IF (DEN.LT.TOL*VMAX) GO TO 180 CON 1200
  V1(J1) = V1(J1)/DEN CON 1210

```


ALGORITHM 518

Incomplete Bessel Function I_0 : The Von Mises Distribution [S14]

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Key Words and Phrases: incomplete Bessel function, von Mises distribution, direction angles, circular statistics, asymptotic normal approximation

CR Categories: 5.5, 5.12

Language: Fortran

DESCRIPTION

The incomplete modified Bessel function $I_0(\theta, \kappa)$ is formally equivalent to the cumulative distribution applied by von Mises [9] to study deviations of atomic weights from integer values, representable as points on the circumference of a circle or as circular directions. This distribution of points on a circle is analogous to the normal or Gaussian distribution of points on a line and has applications to the study of quantal or periodic data, directions of sedimentary bedding, surface fault lines, wildlife movements, etc. (cf. Batschelet [3] and Mardia [7]). The left tail area of this symmetrical distribution is evaluated by this Fortran FUNCTION of the angular deviation θ and the concentration parameter κ , where

$$I_0(\theta, \kappa) = P(x \leq \theta; \kappa) = [2\pi I_0(\kappa)]^{-1} \int_{-\pi}^{\theta} \exp(\kappa \cos x) dx, \quad -\pi \leq \theta < \pi.$$

A compromise between generality and efficiency is achieved by parameterizing the algorithm so that versions corresponding to different levels of accuracy are provided by selecting parameter sets as displayed in Table I.

The method of calculation for small κ involves backwards recursion through a series expansion in terms of modified Bessel functions, while for large κ an auxiliary FUNCTION for the normal probability integral or for the error function $\operatorname{erf}(x)$ is applied to an asymptotic normal approximation to order κ^{-4} .

In the case of small κ the series expansion given by Gumbel et al. [6] can be expressed as

$$\begin{aligned} F(\theta; \kappa) &= \frac{1}{2} + \theta/(2\pi) + \{\pi I_0(\kappa)\}^{-1} \sum_{n=1}^{\infty} n^{-1} I_n(\kappa) \sin(n\theta) \\ &= \frac{1}{2} + \theta/(2\pi) + \pi^{-1} V_1, \end{aligned}$$

where $V_1 = R_1(\kappa) [\sin \theta + R_2(\kappa) [2^{-1} \sin 2\theta + R_3(\kappa) [3^{-1} \sin 3\theta + \dots]]]$ represents a "nested" expression in terms of $R_n(\kappa) = I_n(\kappa)/I_{n-1}(\kappa)$, $n = 1, 2, 3, \dots$,

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Table I. Parameter Values for Several Accuracy Levels

Accuracy in decimal digits	Recurrence over [$a_1 + a_2 \kappa - a_3/(\kappa + a_4)$] terms of nested series				Critical value of κ (CK)	Normalizing approximation (C_1)
	(D)	a_1	a_2	a_3		
6	8	1.0	3.0	1.0	6.5	62.0
7	9	1.0	3.5	0.7	8.0	60.5
8	12	0.8	8	1.0	10.5	56.0
9	15	0.75	16	1.5	15	53.0
10	18	0.7	24	2.0	21	51.7
11	22	0.6	48	3.5	32	50.8
12	28	0.5	100	5.0	50	50.1

the modified Bessel function ratio denoted as r_{n-1} by Amos [2]. Recursive evaluation of the Bessel function ratio is known to be numerically unstable in the direction of increasing n , but backwards recursion generates additional terms of the series expansions in powers of κ , so that numerical stability is achieved for $n = p - 1, \dots, 3, 2, 1$, from the formulation:

$$\sin n\theta = \sin(n+1)\theta \cos \theta - \cos(n+1)\theta \sin \theta$$

$$\cos n\theta = \cos(n+1)\theta \cos \theta + \sin(n+1)\theta \sin \theta$$

$$R_n(\kappa) = [2n/\kappa + R_{n+1}(\kappa)]^{-1} \quad (\text{cf. [2, eq. (2)]})$$

$$V_n = R_n(\kappa) [\sin(n\theta)/n + V_{n+1}].$$

Starting values can be computed directly for $\sin(p\theta)$ and $\cos(p\theta)$, while $R_p(\kappa)$ and V_p may be assigned zero values, provided p is large enough, since $I_p(\kappa)/I_0(\kappa) \leq (\kappa/2)^p/p!$ implies that the truncation error of V_1 will be less than the corresponding truncation error of the convergent series expansion of $\exp(\kappa/2)$. An upper bound on the truncation error was calculated for various combinations of values of κ and p . Curves of $p(\kappa; D)$, interpolated for assigned accuracy levels of D decimal places, were found to be conservatively bounded by relations of the form $p = [a_1 + a_2\kappa - a_3/(\kappa + a_4)]$, where suitable values of coefficients a_1, a_2, \dots , are shown in Table I. Since larger values of κ entail larger values of p and thus more computing time and accumulated roundoff error, an alternative method is desired when κ is larger than some critical value, which we denote by CK .

For large values of κ the asymptotic normality of the von Mises distribution may be exploited in the following three steps. First, the angle θ is transformed to the nearly normally distributed variate $z = b(\kappa) \sin(\theta/2)$, where $b(\kappa) = (2/\pi)^{1/2} \exp(\kappa)/I_0(\kappa)$ is conveniently computed from a continued fraction approximation:

$$b^2(\kappa) \doteq (C - 6 - 54/(C - 26 - 347/(C - C_1)))/6, \quad C = 24\kappa.$$

Values for C_1 , fitted for the minimum value of $\kappa = CK$ are included in Table I. The second step uses an approximation to order κ^{-4} for the asymptotic normalizing series in terms of z [4]:

$$\chi \doteq z - z^3 \left\{ (C - 2z^2 - 16)/3 - \left(z^4 + \frac{7}{4}z^2 + \frac{167}{2} \right) / (C - C_1 - z^2 + 3) \right\}^{-2}$$

in which the same value of C_1 is used in the divisor to reduce the remaining error of order κ^{-5} . The third step calls an auxiliary function, GAUSS (χ), to evaluate the normal integral

$$\Phi(\chi) = (2\pi)^{-1/2} \int_{-\infty}^{\chi} \exp(-t^2/2) dt$$

which, by virtue of the normalizing transforms, provides the required value of the von Mises integral, $F(\theta; \kappa) = \Phi(\chi)$. Alternatively the normal integral may be

evaluated in terms of the error function, $\text{erf}(x)$, by virtue of the relationship, $\Phi(x) = \frac{1}{2}[1 + \text{erf}(x/\sqrt{2})]$, where $\text{erf}(-x) = -\text{erf}(x)$.

The error of the asymptotic approximation decreases rapidly with increasing κ and can be calculated by comparison with sufficiently accurate results from the backwards recursion method. In this way critical values of CK in Table I were obtained such that for $\kappa > CK$ the asymptotic approximation is accurate to at least D decimal digits. Since the number of arithmetic operations for the recursive evaluation increases with increasing κ beyond that required for the asymptotic approximation and the auxiliary function, it is efficient to use backwards recursion for $\kappa \leq CK$ and asymptotic approximation for $\kappa > CK$.

Tests of the algorithm using values from Table I show that resultant probability values are accurate to at least D decimal digits for backward recurrence over $[a_1 + a_2\kappa - a_3/(\kappa + a_4)]$ terms when $\kappa \leq CK$ and for the normal approximations using the corresponding value of C_1 when $\kappa > CK$. For other accuracy levels interpolation should provide useful estimates which may need adjustment to improve efficiency without loss of desired accuracy. The level of accuracy achievable is limited by that of the standard functions for SIN, COS, and SQRT or of the auxiliary function, GAUSS or ERF. For lower-accuracy versions little is lost by use of a linear bound, $p = \kappa + \frac{4}{3}(D - \frac{1}{2})$, but for higher-accuracy versions the tighter bounding curve is more efficient. Should very high accuracy be required, it would be more efficient to utilize high order terms in the asymptotic series [4] to enable reduction of CK .

For real values of θ outside the range $(-\pi, +\pi)$ the result from recursive evaluation lies outside the range $(0,1)$, whereas the normalizing approximation yields results in the required range but lacks the monotonic increasing property appropriate for a cumulative distribution function. It is convenient to extend the domain of definition over all real θ by treating the angle as reduced modulo 2π to the range $(-\pi, +\pi)$, so that results for both small κ and large κ remain consistent with the properties of a probability distribution. With this convention for angles in a circle, the probability over an interval (θ_1, θ_2) , evaluated as $F(\theta_2; \kappa) - F(\theta_1; \kappa)$, can be negative if the interval includes $(2n + 1)\pi$, in which case the correct result is obtained by adding 1. Because accumulated roundoff error can produce results lying just outside $(0,1)$ in extreme cases, such results are replaced by the limit values 0 or 1. It is also convenient to treat negative values of κ as zero so that the result (the cumulative uniform distribution if $\kappa \leq 0$) remains defined for all real κ as well as for all real θ , and thus no error conditions can arise from invalid actual parameter values.

The algorithm is presented in the form of a real-valued FUNCTION VMISES (T, VK). T is the real value of the angle θ in radians, VK is the real value of the concentration parameter κ . This form requires an auxiliary function GAUSS(X) having a single real argument and returning the left tail area of the normal distribution as a real value in the interval $(0,1)$. A suitable function accurate to 10 decimal places can be provided by Algorithm CJ39 [1]. For higher-accuracy levels the somewhat slower algorithm, ACM Algorithm 304 [5] would return the required value with a precision matched to the processor used.

If an auxiliary real function, ERF(X), is supplied instead of GAUSS(X), columns 7 to 50 of each of the last three comment cards marked IF ERF must replace the statement preceding it, thus providing an appropriate alternative version of the function VMISES.

The backwards recurrence method in this algorithm is considerably faster and requires less storage than Algorithm AS 86 [8] which evaluates the von Mises distribution with accuracy up to 8D for θ values extended to the whole real line. The approximation, accurate to 3D, offered for $\kappa > 30$ in the description of Algorithm AS 86, is considerably less effective than the asymptotic approximation in this algorithm.

The performance of this algorithm was tested on an IRIS 80 using a Fortran compiler with single precision to about 7 decimals and double precision to about

16 decimals. The numerical stability of backwards recursive evaluation of $\sin n\theta$ and $\cos n\theta$ was tested in single and double precision for $p \leq 50$, $n = 1$, indicating accuracy loss of at most 2 decimal digits. To check the backwards recurrence method, values of $1 - 2P(\theta; \kappa)$ were computed by the algorithm for $\kappa = 0(0.2)4$ and $\theta = 0^\circ(5^\circ)180^\circ$ and compared with Table 4 of Gumbel et al. [6]. Results corresponded within a unit in the last place, except for two misprints in the published tables; $\Phi_\alpha(\alpha = 10^\circ, \kappa = 1.8) = 0.15740$ should be 0.16740 and $\Phi_\alpha(\alpha = 40^\circ, \kappa = 4.0) = 0.71123$ should be 0.81123. The validity of the asymptotic approximation was tested by comparison with results of extensive recursion for a range of values of θ and κ . Tests with negative and zero values of κ and θ near 0 and 2π and in the range $(-5\pi, +5\pi)$ confirmed performance of the algorithm as specified.

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ALGORITHM

```

FUNCTION VMISES(T, VK)                                VMI 10
C VMISES RETURNS THE LEFT TAIL AREA OF THE VON MISES DISTRIBUTION, VMI 20
C EQUAL TO THE INCOMPLETE MODIFIED BESSEL FUNCTION, OF THE FIRST VMI 30
C KIND AND ZERO-TH ORDER.                               VMI 40
C   T = ANGLE IN RADIANS, TREATED AS DEVIATION FROM ZERO (MEAN) VMI 50
C   REDUCED MODULO 2PI TO THE RANGE (-PI,+PI).          VMI 60
C   VK = CONCENTRATION PARAMETER, KAPPA.  NEGATIVE VALUES ARE VMI 70
C   TREATED AS ZERO.                                    VMI 80
C NEEDS AS AUXILIARY ROUTINE EITHER                     VMI 90
C FUNCTION GAUSS(X), RETURNING THE NORMAL DISTRIBUTION TAIL AREA VMI 100
C   TO THE LEFT OF THE BOUNDING ORDINATE, X,            VMI 110
C OR FUNCTION ERF(X), RETURNING THE ERROR FUNCTION AREA BETWEEN VMI 120
C   ZERO AND THE BOUNDING ORDINATE, X, WHERE ERF(-X) = -ERF(X). VMI 130
C   IF ERF IS USED THEN EACH OF THE LAST 3 COMMENTS (COLUMNS 7 VMI 140
C   TO 50 ONLY) MUST REPLACE THE STATEMENT PRECEDING IT. VMI 150
C   DATA PI /3.1415926535898/, TPI /6.2831853071760/ VMI 160
C CONSTANTS APPROPRIATE FOR 8 DECIMAL DIGIT ACCURACY. VMI 170
C   DATA A1, A2, A3, A4, CK, C1 /12.0,0.8,8.0,1.0,10.5,56.0/ VMI 180
C   Z = VK                                               VMI 190
C CONVERT ANGLE T MODULO 2PI TO RANGE (-PI,+PI).       VMI 200
C   U = AMOD(T+PI,TPI)                                  VMI 210
C   IF (U.LT.0.0) U = U + TPI                           VMI 220
C   Y = U - PI                                           VMI 230
C   IF (Z.GT.CK) GO TO 30                                VMI 240
C   V = 0.0                                              VMI 250
C   IF (Z.LE.0.0) GO TO 20                               VMI 260
C FOR SMALL VK SUM IP TERMS BY BACKWARDS RECURSION. VMI 270
C   IP = Z*A2 - A3/(Z+A4) + A1                          VMI 280
C   P = FLOAT(IP)                                       VMI 290
C   S = SIN(Y)                                           VMI 300
C   C = COS(Y)                                           VMI 310
C   Y = P*Y                                              VMI 320
C   SN = SIN(Y)                                          VMI 330
C   CN = COS(Y)                                          VMI 340
C   R = 0.0                                              VMI 350
C   Z = 2.0/Z                                            VMI 360
C DO 10 N=2,IP                                          VMI 370

```

	P = P - 1.0		VMI 380
	Y = SN		VMI 390
	SN = SN*C - CN*S		VMI 400
	CN = CN*C + Y*S		VMI 410
	R = 1.0/(P*Z+R)		VMI 420
	V = (SN/P+V)*R		VMI 430
	10 CONTINUE		VMI 440
	20 VMISES = (U*0.5+V)/PI		VMI 450
	GO TO 40		VMI 460
C	FOR LARGE VK COMPUTE THE NORMAL APPROXIMATION AND LEFT TAIL.		VMI 470
	30 C = 24.0*Z		VMI 480
	V = C - C1		VMI 490
	R = SQRT((54.0/(347.0/V+26.0-C)-6.0+C)/6.0)		VMI 500
C	R=SQRT((54.0/(347.0/V+26.0-C)-6.0+C)/12.0)	IF ERF	VMI 510
	Z = SIN(Y*0.5)*R		VMI 520
	S = Z*Z		VMI 530
C	S=Z*Z*2.0	IF ERF	VMI 540
	V = V - S + 3.0		VMI 550
	Y = (C-S-S-16.0)/3.0		VMI 560
	Y = ((S+1.75)*S+83.5)/V - Y		VMI 570
	VMISES = GAUSS(Z-S/(Y*Y)*Z)		VMI 580
C	VMISES=ERF(Z-S/(Y*Y)*Z)*0.5+0.5	IF ERF	VMI 590
40	IF (VMISES.LT.0.0) VMISES = 0.0		VMI 600
	IF (VMISES.GT.1.0) VMISES = 1.0		VMI 610
	RETURN		VMI 620
	END		VMI 630

ALGORITHM 519

Three Algorithms for Computing Kolmogorov-Smirnov Probabilities With Arbitrary Boundaries and a Certification of Algorithm 487 [S14]

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Key Words and Phrases: Kolmogorov-Smirnov probabilities, sample distribution function, boundary crossing probability
CR Categories: 5.5
Language: Fortran

DESCRIPTION

Introduction

Let $A(y)$, $B(y)$ be nondecreasing functions on $[0, 1]$ where $B(0) < 0 < A(0)$, $B(1) < 1 < A(1)$. Let $G_n(y)$ be the sample distribution function for n independent random variables uniform on $[0, 1]$. We present three methods for computing $P = \Pr \{B(y) < G_n(y) < A(y), y \in [0, 1]\}$. The special case where $A(y) = y + D$, $B(y) = y - D$ has been treated by ACM Algorithm 487 [4].

Subroutine RAKK. A Generalization of Massey's Method

Massey has given a recursion formula [3; 5, p. 341, eq. (11.7.9)] for computing P when $A(y) = y + k/n$, $B(y) = y - k/n$. This can be generalized to arbitrary boundaries as follows. Let $y(1, j)$, $y(2, j)$ be the level points of $B(y)$, $A(y)$ for the ordinates $(j - 1)/n$ (e.g. $A(y(2, j)) = (j - 1)/n$). If the boundaries are not continuous and strictly increasing, let

$$y(2, j) = \sup\{y \in [0, 1] \mid A(y) \leq (j - 1)/n\}$$

$$y(1, j) = \inf\{y \in [0, 1] \mid B(y) \geq (j - 1)/n\}$$

for those j such that the sup or inf is taken over a nonempty set. Arrange the $y(k, j)$ into a single ordered sequence $y(i)$, $i = 1, \dots, m$, where $y(1) = 0$, $y(m) = 1$ are attached if not already present. Let $jt(i) =$ largest integer j such that $(j - 1)/n < A(y(i) -)$ and $jb(i) =$ smallest integer j such that $(j - 1)/n > B(y(i) +)$, $1 < i < m$. Note that if $jt(m) < n + 1$, then $P = 0$. Then define $jtop(i) = \min(jt(i), n + 1)$, $jbot(i) = \max(jb(i), 1)$, $1 < i < m$; $jbot(1) = jtop(1) = 1$; $jbot(m) = jtop(m) = n + 1$. Then

$$P = \sum_j n! \prod_{k=2}^m \frac{[y(k) - y(k-1)]^{j(k)-j(k-1)}}{j(k)!}$$

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where the summation is over the index set $\{j \equiv (j(1), \dots, j(m)) \mid j_{\text{bot}}(i) \leq j(i) \leq j_{\text{top}}(i)\}$. Note that $j(1) \equiv 1$.

The computer program RAKK sums these multinomial probabilities recursively by computing $W(k, i)$ where

$$W(1, 1) = 1$$

$$W(k, i) = \sum W(j, i-1) [y(i) - y(i-1)]^{k-j} / (k-j)!,$$

$$k = j_{\text{bot}}(i), \dots, j_{\text{top}}(i), \quad i = 1, \dots, m,$$

where the summation is over the index set $\{j \mid k - j \geq 0, j_{\text{bot}}(i-1) \leq j \leq j_{\text{top}}(i-1)\}$. Then $P = W(n+1, m) n!$.

Because machines with small exponent range on floating numbers cannot accommodate the array $W(\cdot, i)$ we have split the array into blocks. Each block contains the values of $W(\cdot, i)$ multiplied by a suitably large conversion factor. The reciprocals of these conversion factors are accumulated in TF by subroutine ACCUM; to keep the accumulated product TF within range it is multiplied by consecutive positive integers, as necessary, with the current value stored in NNN . At the conclusion of the program the value of TF is adjusted so that $NNN = N$ and then the multiplication by $n!$ in the formula $P = W(n+1, m) n!$ is accomplished by a multiplication by TF .

Subroutine DURB. Durbin's Method

The recursion is on $k! \alpha(k)$, $k! \beta(k)$ where α, β are defined in [1, eqs. (36), (37), (40)]. These variables are denoted by $ALPHA(k-1)$, $BETA(k-1)$ in our program, and the indexing runs through $k = 1, \dots, n+1$. Durbin's algorithm computes the probability that $G_n(y)$ crosses a boundary so a subtraction from 1 is necessary. This instruction is flagged in the program listing.

Subroutine EPST. The Epanechnikov, Steck Method

This method [2, p. 15] is numerically unstable for larger values of n because of cancellation of digits in subtractions. Thus we do not include a listing of the computer program. However, we do include the test results for our implementation of this algorithm.

Organization of Programs

All three methods accept as inputs the boundary level points $y(k, j)$ and the sample size n . A preliminary computation of these boundary level points is necessary, and there must be $n+1$ of each. If $A(y) > (j-1)/n$ for all $y \in [0, 1]$, then a negative number must be given for $y(2, j)$ and if $B(y) < (j-1)/n$ for all $y \in [0, 1]$, then a number greater than 1 must be used for $y(1, j)$. As an example, we have included a listing of function PLN which computes P where A, B are (arbitrary) straight lines. As listed, PLN invokes method RAKK; the listing indicates the change necessary to invoke DURB or EPST.

Underflow occurs frequently in our programs and the underflow switch, if any, on the object machine should be set to return zero and the underflow diagnostic should be suppressed.

Test Results, Execution Times, and Precision

The three algorithms were programmed with single precision instructions only and were tested on the following machines: (a) DEC 10, (b) IBM 360-50, (c) CDC 6400. Single precision for these machines is, respectively, 8, 6, and 14 decimal digits and the range on decimal exponents is, respectively, -38 to 38 , -75 to 75 , and -293 to 321 . The fortran compilers used were, respectively, FORTRAN 10/OPT, FORTRAN H(OPT = 2), and FUN. Table I lists the problems tested, where $A(y) = y + D$ and $B(y) = y - D$, and Table II gives the test results. If data is missing, the sample size n is too large for the machine and algorithm involved.

Table I. Problems Tested

Problem	n	D	P (approximate)
1	45	.0527	.001
2	200	.0257	.001
3	45	.1199	.500
4	120	.0742	.500
5	200	.0577	.500
6	45	.2849	.999
7	200	.1368	.999

Table II. Apparent Precision (in significant decimal digits) of Algorithms RAKK, DURB, and EPST

Problem	RAKK			DURB			EPST		
	Machine			Machine			Machine		
	(a)	(b)	(c)	(a)	(b)	(c)	(a)	(b)	(c)
1	7	5	12	4	2	9	7	5	12
2	6	4	11	—	2	9	4	—	10
3	6	4	12	6	5	12	3	2	10
4	6	3	11	—	4	12	—	—	7
5	6	3	11	—	4	11	—	—	4
6	6	3	11	8	6	12	—	—	3
7	6	2	10	2	5	12	—	—	—

Table III. Execution Times (in seconds) of Algorithms RAKK, DURB, and EPST

Problem	RAKK			DURB			EPST		
	Machine			Machine			Machine		
	(a)	(b)	(c)	(a)	(b)	(c)	(a)	(b)	(c)
1	.15	.53	.14	.56	3.08	.48	.03	.12	.02
2	1.73	5.53	1.44	—	71.10	11.21	.30	—	.24
3	.30	1.05	.29	.49	2.80	.41	.05	.32	.05
4	1.77	4.98	1.38	—	22.72	3.51	—	—	.24
5	4.53	11.40	3.22	—	64.52	10.54	—	—	.56
6	.73	1.80	.58	.38	1.92	.29	—	—	.10
7	10.54	28.72	7.24	19.62	53.70	8.84	—	—	—

The precision estimates for machines (a) and (b) are the number of digits which agree with the results of machine (c). The estimates for machine (c) were obtained by comparing the answers for the various methods. Since DURB computes the probability of the complementary event, the final subtraction from 1 accounts for the lesser precision of DURB in problems 1 and 2 and its greater precision in problems 6 and 7.

Execution times for arbitrary boundaries should be comparable to those given in Table III, plus the additional time needed to compute the level points $y(k, j)$.

Recommendations

Although EPST is numerically unstable for larger sample sizes, its execution times are very fast. Thus it would be suitable for numerous computations involving small sample sizes with occasional validations by one of the other routines. DURB is not suitable for large sample sizes on machines with small range on decimal exponents. Execution times for RAKK are faster. Thus RAKK appears to be the preferred choice. This is due largely to the extensive code with its controls against underflow, etc., rather than an intrinsic superiority of the algorithm.

Certification of Algorithm 487

Since Algorithm 487 [4], PKS2, treats a special case handled by Algorithms RAKK, DURB, and EPST, it may be validated by using the data presented in

Table IV. Apparent Precision (in significant decimal digits) of Algorithm PKS2

Machine	Problem Number					
	1	2	3	4	5	6
(a)	—	—	—	—	—	—
(b)	5	—	5	—	—	5
(c)	13+	—	12+	12+	—	13+

Table II. The results for the problems discussed are presented in Table IV. If data are missing the sample size is too large for the machine involved. (A computation of n^n is required by PKS2). Although Pomeranz claims precision of 8 decimal digits on the CDC 6400 [4], this seems rather modest. There is, apparently, a precision of at least 12 decimal digits on the CDC 6400 and a potentially greater precision if the double precision value generated by PKS2 is not returned as a single precision value.

ACKNOWLEDGMENT

These programs are an outgrowth of a seminar in computational probability given during 1974–1975 by Marcel Neuts of Purdue University, who also arranged for computer time on the Purdue machines.

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ALGORITHM

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FUNCTION PLN(N, S2, B2, S1, B1)          PLN  10
C COMPUTES PROB FOR BOUNDARIES OF FORM S*Y+B. N IS          PLN  20
C SAMPLE SIZE, S2,B2 ARE PARAMETERS FOR UPPER BOUNDARY,    PLN  30
C CALLS SUBROUTINE RAKK. TO CALL DURB          PLN  40
C CHANGE AS INDICATED BELOW.          PLN  50
      DIMENSION YY(2,202), S(2), B(2)          PLN  60
C MAX SIZE FOR N IS 200          PLN  70
      S(1) = S1          PLN  80
      S(2) = S2          PLN  90
      B(1) = B1          PLN 100
      B(2) = B2          PLN 110
      IF (N.GT.0) GO TO 10          PLN 120
      PLN = 0.          PLN 130
      RETURN          PLN 140
10 IF (S(1).GE.0.) GO TO 20          PLN 150
   S(1) = 0.          PLN 160
20 IF (S(2).GE.0.) GO TO 30          PLN 170
   B(2) = S(2) + B(2)          PLN 180
   S(2) = 0.          PLN 190
30 N1 = 1 + N          PLN 200
   AN = N          PLN 210
   DO 50 I=1,2          PLN 220
     DO 40 K=1,N1          PLN 230
       C = FLOAT(K-1)/AN - B(I)          PLN 240
       YY(I,K) = -.5          PLN 250
       IF (C.LE.0.) GO TO 40          PLN 260
       YY(I,K) = 1.1          PLN 270
       IF (S(I).GT.C) YY(I,K) = C/S(I)          PLN 280
40 CONTINUE          PLN 290
50 CONTINUE          PLN 300
   CALL RAKK(N, YY, ANS)          PLN 310

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COLLECTED ALGORITHMS (cont.)

519-P 5- 0

C TO CALL DURB CHANGE PRECEDING STATEMENT TO	PLN 320
C CALL DURB (N,YY,ANS)	PLN 330
PLN = ANS	PLN 340
RETURN	PLN 350
END	PLN 360
SUBROUTINE RAKK(N, YY, ANS)	RAK 10
C N IS SAMPLE SIZE. YY(2,.),YY(1,.)	RAK 20
C CONTAIN, RESPECTIVELY, UPPER AND LOWER BOUNDARY	RAK 30
C LEVEL POINTS CORRESPONDING TO ϕ , 1/N, 2/N, ..., 1. MUST	RAK 40
C BE N+1 OF EACH. ROUTINE ASSUMES THEY ARE NONDECREASING.	RAK 50
DIMENSION W(201,2), YY(2,202), KBLOC(41), FRNG(41,2), SSS(41)	RAK 60
DATA RNG, RRNG, SQRNG, SQRRNG, MAXNB /1.E292,1.E-292,1.E146,	RAK 70
* 1.E-146,40/	RAK 80
C RNG AND RRNG ARE MACHINE DEPENDENT CONSTANTS SELECTED	RAK 90
C SO RRNG=1./RNG G.E. SMALLEST POSITIVE FLOATING	RAK 100
C NUMBER AND RNG L.E. LARGEST POSITIVE FLOATING	RAK 110
C NUMBER OF MACHINE. SQRNG AND SQRRNG ARE SQUARE ROOTS	RAK 120
C OF RNG AND RRNG.	RAK 130
C MAXNB IS THE MAX NUMBER OF BLOCKS INTO WHICH	RAK 140
C W(.,.) WILL BE PARTITIONED. IF MAXNB IS	RAK 150
C INCREASED REVISE DIMENSION STATEMENT ACCORDINGLY.	RAK 160
N1 = N + 1	RAK 170
TF = 1.	RAK 180
NNN = 1	RAK 190
IEXT = 1	RAK 200
M2 = 1	RAK 210
M1 = 2	RAK 220
W(1,M2) = SQRNG	RAK 230
NBLOC = 1	RAK 240
KBLOC(1) = 1	RAK 250
KBLOC(2) = 2	RAK 260
FRNG(1,M2) = 1.	RAK 270
ANS = ϕ .	RAK 280
IF (N1.LE.1) RETURN	RAK 290
YY(2,N1+1) = 1.	RAK 300
Y1 = ϕ .	RAK 310
IF ((YY(1,1).LE. ϕ .) .OR. (YY(2,N1).GE.1.) .OR.	RAK 320
* (YY(1,N1).LE.1.) .OR. (YY(2,1).GE. ϕ .) RETURN	RAK 330
JBOTP = 1	RAK 340
JTOPP = 1	RAK 350
JTOP1 = 1	RAK 360
C ENTRY POINT FOR SUBSEQUENT ITERATIONS.	RAK 370
10 JBOT ϕ = JBOTP	RAK 380
JTOP ϕ = JTOP1	RAK 390
MT = M1	RAK 400
M1 = M2	RAK 410
M2 = MT	RAK 420
Y ϕ = Y1	RAK 430
C COMPUTE SUMMATION INDICES, JBOTP,JTOPP FOR NEXT	RAK 440
C ITERATION.	RAK 450
20 IF (YY(2,JTOPP+1).GT.Y ϕ) GO TO 30	RAK 460
JTOPP = JTOPP + 1	RAK 470
GO TO 20	RAK 480
30 Y1 = AMIN1(YY(2,JTOPP+1),YY(1,JBOTP))	RAK 490
IF (Y1.GE.1.) GO TO 50	RAK 500
40 IF (YY(1,JBOTP).GT.Y1) GO TO 60	RAK 510
JBOTP = JBOTP + 1	RAK 520
GO TO 40	RAK 530
50 Y1 = 1.	RAK 540
C SET EXIT FLAG. NEXT ITERATION IS FINAL ONE.	RAK 550
IEXT = 2	RAK 560
JBOTP = N1	RAK 570
60 IF (JBOTP.GT.JTOPP) RETURN	RAK 580
C RETURN WITH ANS= ϕ SINCE NO PATHS BETWEEN BOUNDARIES.	RAK 590
JTOP1 = JTOPP	RAK 600
P = Y1 - Y ϕ	RAK 610
KBLOC(NBLOC+1) = JTOP1 + 1	RAK 620
DO 70 L=JBOT ϕ ,JTOP1	RAK 630
W(L,M2) = ϕ .	RAK 640
70 CONTINUE	RAK 650
L ϕ = KBLOC(NBLOC)	RAK 660
L1 = JTOP ϕ	RAK 670
DO 80 L=L ϕ ,L1	RAK 680

IF (W(L,M1).GT.SQRRNG) GO TO 80	RAK 690
JTOP0 = L - 1	RAK 700
GO TO 90	RAK 710
80 CONTINUE	RAK 720
90 IF (JBOT0.LT.KBLOC(2)) GO TO 110	RAK 730
C THE NEXT FEW STATEMENTS ELIMINATE BLOCKS NO	RAK 740
C LONGER USED.	RAK 750
DO 100 I=1,NBLOC	RAK 760
FRNG(I,M1) = FRNG(I+1,M1)	RAK 770
KBLOC(I) = KBLOC(I+1)	RAK 780
100 CONTINUE	RAK 790
NBLOC = NBLOC - 1	RAK 800
CALL ACCUM(RRNG/FRNG(1,M1), TF, NNN)	RAK 810
GO TO 90	RAK 820
110 KBLOC(1) = JBOT0	RAK 830
MBAV = (MAXNB-NBLOC+1)/2	RAK 840
IF (P.LE.(1./FLOAT(4*N1))) MBAV = 2	RAK 850
MBAV = MIN0(MAXNB,NBLOC+MBAV)	RAK 860
C MBAV IS UPPER LIMIT ON NUMBER OF BLOCKS DURING NEXT	RAK 870
C ITERATION. MBAV IS SET TO LIMIT NUMBER OF NEW	RAK 880
C BLOCKS WHEN P IS SMALL.	RAK 890
CONVI = 1.	RAK 900
I1 = 1	RAK 910
J1 = 1	RAK 920
120 DO 400 I=I1,NBLOC	RAK 930
KBLI = KBLOC(I)	RAK 940
IF (I.EQ.I1) GO TO 170	RAK 950
FRNG(I,M2) = AMINI(FRNG(I,M1), 1./((SQRRNG*W(KBLI-1,M2))))	RAK 960
CONVI = (CONVI/FRNG(I-J1+1,M1))*FRNG(I,M2)	RAK 970
IF (J1.EQ.1) GO TO 170	RAK 980
K1 = KBLOC(I) - KBLOC(I-J1+2) + 1	RAK 990
K0 = KBLOC(I-1) - KBLOC(I-J1+1) + 1	RAK 1000
IF (K1-K0) 150, 170, 130	RAK 1010
130 K0P1 = K0 + 1	RAK 1020
DO 140 K=K0P1,K1	RAK 1030
CONVI = CONVI*P/FLOAT(K)	RAK 1040
140 CONTINUE	RAK 1050
GO TO 170	RAK 1060
150 K1P1 = K1 + 1	RAK 1070
DO 160 K=K1P1,K0	RAK 1080
CONVI = CONVI*FLOAT(K)/P	RAK 1090
160 CONTINUE	RAK 1100
170 IF (CONVI.GT.RRNG) GO TO 210	RAK 1110
DO 190 J=J1,J11	RAK 1120
C CHANGE UPPER LIMIT IF EARLY EXIT FROM PREVIOUS LOOP	RAK 1130
CONVI = FRNG(I,M2)*RNG	RAK 1140
K0 = MAX0(1,KBLOC(I-1)-MIN0(JTOP0+1,KBLOC(I-J+1))+2)	RAK 1150
K1 = KBLOC(I) - MIN0(JTOP0+1,KBLOC(I-J+1)) + 1	RAK 1160
DO 180 K=K0,K1	RAK 1170
CONVI = CONVI*P/FLOAT(K)	RAK 1180
180 CONTINUE	RAK 1190
CONVI = SSS(J)*CONVI	RAK 1200
IF (CONVI.GT.RRNG) GO TO 200	RAK 1210
190 CONTINUE	RAK 1220
GO TO (10, 440), IEXT	RAK 1230
200 J1 = J + 1	RAK 1240
210 DO 390 J=J1,I	RAK 1250
LFG = 1	RAK 1260
C LFG=2 IF NEXT ITERATION ON J BECOMES NECESSARY.	RAK 1270
LBI = KBLOC(I)	RAK 1280
K01 = KBLOC(I) - KBLOC(I-J+1) + 1	RAK 1290
S = CONVI	RAK 1300
IF (J.EQ.J1) GO TO 240	RAK 1310
IF (TMLT) 220, 400, 230	RAK 1320
220 S = -TMLT*FRNG(I-J+2,M1)	RAK 1330
IF (S.GT.RRNG) GO TO 240	RAK 1340
230 S = TMLT*(FRNG(I-J+2,M1)*RNG)	RAK 1350
240 K0 = MAX0(KBLOC(I)-MIN0(KBLOC(I-J+2)-1,JTOP0)+1,1)	RAK 1360
K1 = KBLOC(I+1) - KBLOC(I-J+1)	RAK 1370
TMLT = 0.	RAK 1380
RRNGS = 1.	RAK 1390
IF (S.LT.0.) RRNGS = -RRNG	RAK 1400
SSS(J) = S	RAK 1410
DO 370 K=K0,K1	RAK 1420
LDF = 1 - K	RAK 1430

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      LB = MAX0(KBLOC(I),KBLOC(I-J+1))-LDF,LB1)
      LT = MIN0(KBLOC(I+1)-1,MIN0(JTOP0,KBLOC(I-J+2)-1))-LDF)
      IF (LB.GT.LB1) LFG = 2
      IF (LB.GT.LT) GO TO 270
      DO 260 L=LB,LT
      *   IF (W(L,M2).EQ.((W(L+LDF,M1)*RRNGS)*S+W(L,M2))) GO TO
        250
        LB1 = L
C WHEN THE SUMMANDS BECOME SO SMALL THAT THE VALUE OF
C W(L,M2) IS UNCHANGED LB1 IS INCREASED. THIS PREVENTS
C FURTHER ADDITIONS TO SUCH W(L,M2).
      GO TO 290
250     IF (W(L,M2).GT.RRNG) GO TO 260
      IF (K-K01) 340, 340, 280
260     CONTINUE
270     LB1 = LT + 1
      IF ((LB1.NE.KBLOC(I+1)) .OR. (K.LE.K01)) GO TO 340
280     J11 = J
      IF (W(KBLI,M2).LE.RRNG) LFG = 2
      GO TO (400, 380), LFG
290     IF (S) 300, 380, 320
300     DO 310 L=LB1,LT
      W(L,M2) = (W(L+LDF,M1)*RRNGS)*S + W(L,M2)
310     CONTINUE
      GO TO 340
320     DO 330 L=LB1,LT
      W(L,M2) = W(L+LDF,M1)*S + W(L,M2)
330     CONTINUE
340     T = S
      S = S*P/FLOAT(K)
      IF (S) 360, 350, 360
350     IF (T.LE.0.) GO TO 380
      S = -(T*RRNG)*P/FLOAT(K)
      RRNGS = -RRNG
C THE NEGATIVE BIT IN S,TMLT IS USED TO
C INDICATE ITS VALUE IS RNG TIMES THE USUAL VALUE.
360     IF (K.EQ.K01) TMLT = S
C THE ITERATION ON K MUST CONTINUE UNTIL TMLT IS
C SET. TMLT USED IN NEXT ITERATION ON J.
370     CONTINUE
380     J11 = J
390     CONTINUE
400     CONTINUE
C CREATE NEW BLOCK IF NECESSARY.
      IF (NBLOC.GE.MBAV) GO TO (10, 440), IEXT
      L1 = KBLOC(NBLOC+1) - KBLOC(NBLOC)
      DO 410 L=1,L1
      LL = KBLOC(NBLOC+1) - L
      IF (W(LL,M2).GE.SQRRNG) GO TO 420
410     CONTINUE
420     IF (LL.EQ.JTOP1) GO TO (10, 440), IEXT
      KBLOC(NBLOC+2) = KBLOC(NBLOC+1)
      KBLOC(NBLOC+1) = LL + 1
      NBLOC = NBLOC + 1
      L0 = KBLOC(NBLOC)
      DO 430 L=L0,JTOP1
      W(L,M2) = 0.
430     CONTINUE
      I1 = NBLOC
      FRNG(NBLOC,M2) = 1.
      CONVI = 0.
      GO TO 120
440     CALL ACCUM(W(N1,M2), TF, NNN)
      CALL ACCUM(SQRRNG, TF, NNN)
      IF (NBLOC.EQ.1) GO TO 460
      DO 450 I=2,NBLOC
      CALL ACCUM(RRNG/FRNG(I,M2), TF, NNN)
450     CONTINUE
460     ANS = TF
      IF (N-NNN) 470, 490, 500
470     DO 480 J=N1,NNN
      ANS = ANS/FLOAT(J)
480     CONTINUE
490     RETURN
500     NNN = NNN + 1

```

RAK 1440
 RAK 1450
 RAK 1460
 RAK 1470
 RAK 1480
 RAK 1490
 RAK 1500
 RAK 1510
 RAK 1520
 RAK 1530
 RAK 1540
 RAK 1550
 RAK 1560
 RAK 1570
 RAK 1580
 RAK 1590
 RAK 1600
 RAK 1610
 RAK 1620
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 RAK 2010
 RAK 2020
 RAK 2030
 RAK 2040
 RAK 2050
 RAK 2060
 RAK 2070
 RAK 2080
 RAK 2090
 RAK 2100
 RAK 2110
 RAK 2120
 RAK 2130
 RAK 2140
 RAK 2150
 RAK 2160
 RAK 2170
 RAK 2180

DO 510 J=NNN,N	RAK 2190
ANS = ANS*FLOAT(J)	RAK 2200
510 CONTINUE	RAK 2210
RETURN	RAK 2220
END	RAK 2230
SUBROUTINE ACCUM(F, TF, NNN)	ACC 10
C THIS SUBROUTINE ANCILLARY TO RAKK AND ACCUMULATES THE	ACC 20
C MULTIPLIERS USED TO KEEP W(.,.) IN RANGE.	ACC 30
TF = TF*F	ACC 40
10 IF ((TF.GT.1.) .OR. (TF.EQ.0.)) RETURN	ACC 50
NNN = NNN + 1	ACC 60
TF = TF*FLOAT(NNN)	ACC 70
GO TO 10	ACC 80
END	ACC 90
SUBROUTINE DURB(N, YY, ANS)	DUR 10
C DURBINS METHOD. J.APPL.PROB., 8(1971), PP431-453,	DUR 20
C FORMULAS (36),(37),(40).	DUR 30
C N IS SAMPLE SIZE. YY(2,.) ,YY(1,.)	DUR 40
C CONTAIN, RESPECTIVELY, UPPER AND LOWER BOUNDARY	DUR 50
C LEVEL POINTS CORRESPONDING TO 0,1/N,2/N,...,1. MUST	DUR 60
C BE N+1 OF EACH. ROUTINE ASSUMES THEY ARE NONDECREASING.	DUR 70
DIMENSION ALPHA(201), BETA(201), YY(2,202)	DUR 80
C DIMENSION STATEMENT PERMITS MAX SAMPLE SIZE OF 200	DUR 90
N1 = N + 1	DUR 100
ANS = 0.	DUR 110
IF ((YY(1,1).LE.0.) .OR. (YY(2,N1).GE.1.) .OR.	DUR 120
* (YY(1,N1).LE.1.) .OR. (YY(2,1).GE.0.)) RETURN	DUR 130
IP = 1	DUR 140
DO 20 I=1,N1	DUR 150
IF (YY(2,I).LT.0.) GO TO 10	DUR 160
IP = I	DUR 170
GO TO 30	DUR 180
10 YY(2,I) = 0.	DUR 190
20 CONTINUE	DUR 200
30 DO 40 I=1,N1	DUR 210
IF (YY(1,I).LE.1.) GO TO 40	DUR 220
YY(1,I) = 1.	DUR 230
40 CONTINUE	DUR 240
ALPHA(1) = 0.	DUR 250
BETA(1) = 1.	DUR 260
DO 120 K=2,N1	DUR 270
ALPHA(K) = YY(2,K)**(K-1)	DUR 280
IF (IP.GT.(K-1)) GO TO 60	DUR 290
COMB = 1.	DUR 300
I1 = K - IP	DUR 310
DO 50 I=1,I1	DUR 320
COMB = COMB*FLOAT(K-I)/FLOAT(I)	DUR 330
C COMB IS BINOMIAL COEFFICIENT K-1 OVER I-1.	DUR 340
ALPHA(K) = ALPHA(K) - ALPHA(K-I)*(YY(2,K)-YY(2,K-I))**I*	DUR 350
* COMB	DUR 360
50 CONTINUE	DUR 370
60 COMB = 1.	DUR 380
DO 70 J=1,N1	DUR 390
IF (YY(2,K).LE.YY(1,J)) GO TO 80	DUR 400
ALPHA(K) = ALPHA(K) - BETA(J)*(YY(2,K)-YY(1,J))**(K-J)*	DUR 410
* COMB	DUR 420
C COMB IS BINOMIAL COEFFICIENT K-1 OVER J-1.	DUR 430
COMB = COMB*FLOAT(K-J)/FLOAT(J)	DUR 440
70 CONTINUE	DUR 450
80 BETA(K) = YY(1,K)**(K-1)	DUR 460
IF (IP.GT.K) GO TO 100	DUR 470
COMB = 1.	DUR 480
I1 = K - IP + 1	DUR 490
DO 90 I=1,I1	DUR 500
BETA(K) = BETA(K) - ALPHA(K-I+1)*(YY(1,K)-YY(2,K-I+1))**	DUR 510
* (I-1)*COMB	DUR 520
C COMB IS BINOMIAL COEFFICIENT K-1 OVER I-1.	DUR 530
COMB = COMB*FLOAT(K-I)/FLOAT(I)	DUR 540
90 CONTINUE	DUR 550
COMB = 1.	DUR 560

100	KMI = K - 1	DUR	570
	DO 110 J=1,KMI	DUR	580
	IF (YY(1,J).EQ.1.) GO TO 120	DUR	590
C	COMB IS BINOMIAL COEFFICIENT K-1 OVER J-1.	DUR	600
	BETA(K) = BETA(K) - BETA(J)*(YY(1,K)-YY(1,J))**(K-J)*COMB	DUR	610
	COMB = COMB*FLOAT(K-J)/FLOAT(J)	DUR	620
110	CONTINUE	DUR	630
120	CONTINUE	DUR	640
	COMB = 1.	DUR	650
	I1 = N1 - IP + 1	DUR	660
	DO 130 I=1,I1	DUR	670
	ANS = ANS + (1.-YY(2,N1-I+1))**(I-1)*ALPHA(N1-I+1)*COMB	DUR	680
	COMB = COMB*FLOAT(N1-I)/FLOAT(I)	DUR	690
130	CONTINUE	DUR	700
	COMB = 1.	DUR	710
	DO 140 J=1,N1	DUR	720
	IF (YY(1,J).EQ.1.) GO TO 150	DUR	730
	ANS = ANS + (1.-YY(1,J))**(N1-J)*BETA(J)*COMB	DUR	740
C	COMB IS BINOMIAL COEFFICIENT N OVER J-1.	DUR	750
	COMB = COMB*FLOAT(N1-J)/FLOAT(J)	DUR	760
140	CONTINUE	DUR	770
C	FOR PROB OF COMPLEMENTARY EVENT CHANGE NEXT STATEMENT TO	DUR	780
C	150 CONTINUE	DUR	790
150	ANS = 1. - ANS	DUR	800
	RETURN	DUR	810
	END	DUR	820

ALGORITHM 520

An Automatic Revised Simplex Method for Constrained Resource Network Scheduling [H]

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Key Words and Phrases: resource allocation, activity networks, linear programming

CR Categories: 3.31, 4.32, 5.41, 8.3

Language: Fortran

DESCRIPTION

Purpose

Subroutine ARSME solves a resource constrained, network scheduling problem for the case in which activities may be arbitrarily interrupted and restarted later with no increase in activity duration. The number of resource types is not a limiting factor in our procedure. The amount of any one resource available at any moment is constant. We shall use the "activity-on-arc" network representation, under the commonly imposed assumption that the network contains no directed cycles and has only one "beginning" and only one "terminal" node (event). It is further assumed that the network nodes (events) are ordered in such a way that node i precedes node j , if $i < j$. Such an ordering is always possible and it induces an ordering among the arcs (activities).

Optimal approaches to resource constrained, network scheduling problems where activities can require more than one resource type are presented in [1, 4]. Both methods assume integer durations of activities, and the method presented in [1] divides activity durations into unit intervals. Both methods can handle networks with up to about 30 activities and 3 resource types.

Subroutine ARSME is constructed in such a way that its storage requirements are minimal, a fact which permits the solution of problems for very large networks with many resource types. Moreover, an optimal solution can be obtained in a shorter time when relatively smaller amounts of the resources are available than when resources are less limited.

Let the number of activities be equal to M and the number of resource types be equal to RT . For activity j ($j = 1, 2, \dots, M$) and resource k ($k = 1, 2, \dots, RT$) the following variables are given: activity duration d_j ; activity resource requirement R_{jk} , i.e. the amount of resource k required by activity j ; resource limit RL_k , i.e. the amount of resource k available at any moment.

The solution of our problem consists of the assignment of time intervals to resource feasible sets of activities (i.e. activities which can be performed simul-

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taneously and not violate resource limits), in order that the completion time of the project be minimal. For this reason, it is necessary to generate all such sets and to solve the ensuing linear programming (LP) representation of the problem. LP variables are the durations assigned to these activity sets. The optimal durations obtained are ordered according to the network structure. In similar fashion we obtain the time intervals making up the optimal schedule.

As input data for subroutine ARSME, it is necessary to provide the following data:

M	number of activities in the network;
$M1$	dimension limit equal to $M + 1$;
$M2$	dimension limit equal to $2M + 1$;
ME	dimension limit of vector E (see the discussion below);
ND	number of nodes;
RT	number of resource types;
NI	declared limit of the number of iterations, usually taken as $2M$ (the subroutine ends computations if the optimum solution is not found in NI iterations);
$NS(M2)$	network structure vector containing the consecutive activities as ordered pairs of nodes; no initial value is required for element $NS(M2)$;
$X(M1)$	vector of activity durations;
$RL(RT)$	vector of resource limits;
$R(M, RT)$	matrix of resource requirements.

The value of ME is taken to be $(M1 \times NI)$ if $M1 \times NI$ does not exceed the available amount of primary storage; otherwise, it is taken as the available limit of primary storage.

When the optimal solution is obtained, the optimal schedule (i.e. the sets of activities together with corresponding time intervals) are printed on the line printer. On the other hand, if the optimal solution is not reached in the declared number of iterations NI , then the message "number of iterations greater than NI " is printed. In this case, the computations should be repeated for a greater value of NI .

As an example let us consider the network shown in Figure 1. Input data are: $M = 5$; $M1 = 6$; $M2 = 11$; $ME = 60$; $ND = 4$; $RT = 3$; $NI = 10$; $NS = [1, 2, 1, 3, 2, 3, 2, 4, 3, 4]$; $X = [1.0, 2.0, 3.0, 2.0, 2.0]$; $RL = [5.0, 5.0, 3.0]$;

$$R = \begin{pmatrix} 2.0 & 2.0 & 1.0 \\ 0.0 & 2.0 & 1.0 \\ 2.0 & 1.0 & 3.0 \\ 3.0 & 3.0 & 3.0 \\ 1.0 & 1.0 & 0.0 \end{pmatrix}.$$

As an optimal schedule, ARSME yields

time intervals	activities				
	1	2	3	4	5
0.0-1.0	1	1	0	0	0
1.0-2.0	0	1	0	0	0
2.0-5.0	0	0	1	0	0
5.0-7.0	0	0	0	1	1
7.0-7.0	0	0	0	0	1

with $T_{\min} = 7$. This means that, in the optimal schedule, activities 1 and 2 should be performed simultaneously in the time interval $(0.0, 1.0)$, activity 2 in the time interval $(1.0, 2.0)$, etc.

Method

Subroutine ARSME solves an LP problem which is formulated as follows.

Let S_T denote the set of all activities which may be performed between the oc-

currence of node I and $I + 1$. Such sets will be called *main sets*. Let us number from 1 to N the resource feasible sets, i.e. those subsets of the main sets and those main sets for which the resource requirements do not exceed any resource limits.

Now, let Q_j denote the set of all numbers of resource feasible sets in which activity j may be performed, and let x_i denote the duration of set i . Thus the LP problem is obtained:

minimize

$$T = \sum_{i=1}^N x_i$$

subject to

$$\sum_{i \in Q_j} x_i = d_j \quad (j = 1, 2, \dots, M) \quad (1)$$

or in matrix notation,

$$\mathbf{A} \mathbf{x} = \mathbf{d} \quad (2)$$

where \mathbf{A} is the matrix of coefficients

$$a_{ji} = \begin{cases} 1 & \text{if } i \in Q_j; \\ 0 & \text{otherwise.} \end{cases}$$

Obviously, columns of matrix \mathbf{A} correspond to resource feasible sets.

When applying the simplex method directly to this problem, one must somehow store matrix \mathbf{A} , which is a very large matrix for even small problems. This greatly constrains the size of the problems which can be solved.

Subroutine ARSME eliminates the above storage problem by utilizing the revised simplex method [2] in which elements of matrix \mathbf{A} are not transformed. This method is computationally profitable for sequencing problems, since in practical ones, the number of variables is greater than three times the number of constraints [5]. The specific version used is based upon the product form of the inverse [2].

The first basic feasible solution to the problem is known in advance. It consists of M , 1-element sets (consecutive performance of all activities).

Taking advantage of the special structure of the problem, a subroutine GEN has been developed which automatically generates the consecutive columns of matrix \mathbf{A} (i.e. the resource feasible sets). An important feature of this subroutine is that at every moment only one main set and only one of its resource feasible subsets (which is actually needed in the simplex procedure) are stored. This allows one to solve problems of larger dimension than otherwise possible and of course radically simplifies the preparation of input data.

Subroutine GEN generates main sets S_1, S_2, \dots, S_{ND} using the network structure vector. When a single main set has been generated, the generation of its resource feasible subsets follows. For this purpose, so-called *primary subsets* are constructed, from which resource feasible subsets are created by adding activities from the main set. The first primary subset contains all the activities of the main set which also composed previously generated main sets, as well as one new activity. New primary subsets are created either when the last activity from the main set has been added, or when the generated subset is found to be unfeasible. In the first case, the last activity of the last generated subset is rejected and the penultimate one is replaced by the next from the main set. In the second case, the last activity of the unfeasible subset is replaced by the next from the main set. The number of sets in which resource feasibility is checked is always minimal and is usually much smaller than the total number of sets of activities which may be performed simultaneously. This follows from the construction of the primary subsets and from the fact that activities in the main sets are ordered by the subroutine ORDSNE in a unique fashion.

After generating each resource feasible set ARSME, together with subrou-

tine K FIND, checks, according to the simplex procedure, whether the set is good enough to be in the optimal solution. This procedure is repeated until the optimal solution is obtained. Then subroutine ORDSNE orders the sets composing the optimal solution, and subroutine PRISET prints out the results.

The parameters of ARSME, which are associated with the applied version of the simplex procedure, are as follows:

- U, D vector used for finding vectors introduced into, and eliminated from, the basis, respectively;
- E vector of the inverse;
- LO vector of the indices of vectors eliminated from the basis in successive iterations;
- IBV vector of indices of the basis vectors.

(Other parameters are subsidiary ones).

Timing Tests

A comparison of the computation time required by subroutine ARSME and other methods for solving the constrained resource network scheduling problem is not very helpful because of the different assumptions made in these methods. However, as an example, a comparison of computation times is shown in Table I for the two

Table I

Problem	Resource limits RL_k	Davis method		Patterson method		ARSME	
		Optimal schedule length	Time to optimal solution*	Optimal schedule length	Time to optimal solution†	Optimal schedule length§	Time to optimal solution‡
1. Davis and Heidorn	5, 5, 3	7		7	1.57	7	0.81
2. Moodie and Mandeville	5	8	2.00	8	1.35	8	1.20
3. Moodie and Mandeville	4	11		11	3.78	10.5	1.14

* IBM 7094 CPU time in seconds.
 † IBM 360/67 CPU time in seconds.
 ‡ ICL 1900 CPU time in seconds.
 § Because of the provision for job splitting in our procedure, optimal schedule lengths may differ from those reported by Davis and by Patterson.

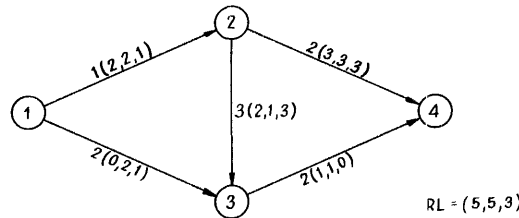


Fig. 1. Network used in problem formulated by Davis and Heidorn [1]

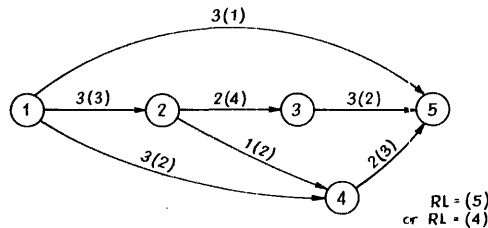


Fig. 2. Network used in problem formulated by Moodie and Mandeville [3]

simple networks described in [1, 3] and shown in Figures 1 and 2. One may notice that, for problem 3, ARSME finds a solution with a lower value of the objective function, because of the assumption of the arbitrary splitting of activities.

All tests have been made on an Odra-1305 computer (license ICL-1900). Computation times for networks of about 20 activities vary from about 15 seconds to 3 minutes; for networks of about 50 activities, from 5 minutes to 25 minutes; and for networks of about 100 activities, from 20 minutes to 90 minutes. Variance in computation times for a given size of network appears to be primarily a function of the number of network nodes and the values of resource limits.

ACKNOWLEDGMENTS

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ALGORITHM

[Only that portion of the listing which is the introductory comment section explaining subroutine ARSME is printed here. The complete listing, including subroutines GEN, KFINd, PRISet, and ORDSNE, is available from the ACM Algorithms Distribution Service (see inside back cover for order form), or may be found in "Collected Algorithms from ACM."]

```

      SUBROUTINE ARSME(M, M1, M2, ME, ND, RT, NI, NS, X, R, RL,      ARS  10
      * IBV, LO, E, U, D)                                         ARS  20
C PURPOSE                                                         ARS  30
C TO SOLVE A MULTIPLE-RESOURCE NETWORK SCHEDULING PROBLEM -PREEMPTIVE ARS  40
C CASE BY THE REVISED SIMPLEX METHOD WITH THE PRODUCT FORM OF THE ARS  50
C INVERSE. *ACTIVITIES-ON-ARCS* NETWORK REPRESENTATION IS USED. ARS  60
C THIS IS THE PRIMARY SUBROUTINE AND IT COORDINATES THE SPECIAL ARS  70
C PURPOSE SUBROUTINES GEN,KFINd,PRISet,ORDSNE.                  ARS  80
C THE SUBROUTINE USES A LINE-PRINTER AS PROGRAMMING UNIT 2 AND ARS  90
C MAGNETIC TAPES AS PROGRAMMING UNIT 3.                          ARS 100
C THE MAGNETIC TAPES ARE USED ONLY IF THE INVERSE EXCEEDS THE ARS 110
C DIMENSION LIMIT OF VECTOR E.                                   ARS 120
C DESCRIPTION OF PARAMETERS                                       ARS 130
C M   NUMBER OF ACTIVITIES                                       ARS 140
C M1  DIMENSION LIMIT OF VECTORS IBV,U,D AND X, EQUAL TO M+1   ARS 150
C M2  DIMENSION LIMIT OF VECTOR NS, EQUAL TO M+M+1              ARS 160
C ME  DIMENSION LIMIT OF VECTOR E                               ARS 170
C ND  NUMBER OF NODES                                           ARS 180
C RT  NUMBER OF RESOURCE TYPES                                   ARS 190
C NI  DECLARED LIMIT OF THE NUMBER OF ITERATIONS ( ABOUT 2*M ) ARS 200
C NS  NETWORK STRUCTURE VECTOR CONTAINING THE CONSECUTIVE ACTIVITIES ARS 210
C     AS ORDERED PAIRS OF NODES.                                 ARS 220
C     NO INITIAL VALUE IS REQUIRED FOR ELEMENT NS(M2).           ARS 230
C X   VECTOR OF ACTIVITY DURATIONS                               ARS 240
C R   MATRIX OF RESOURCE REQUIREMENTS                            ARS 250
C RL  VECTOR OF RESOURCE LIMITS                                  ARS 260
C E   VECTOR OF THE INVERSE                                      ARS 270
C U,D VECTORS USED FOR FINDING VECTORS INTRODUCED INTO AND      ARS 280
C     ELIMINATED FROM THE BASIS RESPECTIVELY                    ARS 290
C LO  VECTOR OF THE INDICES OF VECTORS ELIMINATED FROM THE BASIS IN ARS 300
C     SUCCESSIVE ITERATIONS                                      ARS 310
C IBV VECTOR OF INDICES OF THE BASIC VECTORS                     ARS 320
C AS INPUT DATA PROVIDE M,M1,M2,ME,ND,RT,NI,NS,X,R,RL.        ARS 330

```

```

      INTEGER RT, RT1
      LOGICAL PRE, PROC
      DIMENSION NS(M2), IBV(M1), LO(NI), E(ME), U(M1), D(M1),
      * X(M1), R(M,RT), RL(RT)
C FOR SOME COMPILERS OF FORTRAN THE DIMENSION LIMIT OF A MATRIX MUST
C BE GREATER THAN 1, THEN IF RT IS TO BE EQUAL TO 1, MAKE RT=2 AND
C RL(2)=0
      PROC = .FALSE.
      PRE = .TRUE.
      RT1 = RT
      IF (RL(RT).EQ.0.) RT1 = RT - 1
      LG = 0
      IT = 0
      JK = 0
      IH2 = 0
      IH = 1
      JP = 1
      ND = ND - 1
      NS(M2) = NS(M2-1) + 1
      IH1 = ME/M1
      RR = 0.
      DO 10 I=1,M
         RR = RR + X(I)
10 CONTINUE
      X(M1) = -RR
C CHECK IF THE RESOURCE REQUIREMENTS FOR ANY TYPE OF RESOURCE ARE
C EQUAL TO EACH OTHER. IF SO, CALCULATE IP IT IS THE MAXIMUM NUMBER
C OF ACTIVITIES WHICH MAY BE PERFORMED SIMULTANEOUSLY.
      IP = 99999
      DO 30 J=1,RT1
         P = R(1,J)
         DO 20 I=2,M
            IF (P.NE.R(I,J)) GO TO 40
20 CONTINUE
            IF (P.LT.0.000001) GO TO 30
            I = (RL(J)+.00001)/P
            IF (IP.GT.I) IP = I
30 CONTINUE
            IF (IP.NE.99999) PROC = .TRUE.
40 D(M1) = 0.
C CALCULATE VECTOR U
      DO 50 I=1,M
         U(I) = 0.
50 CONTINUE
      U(M1) = 1.
      MF = IH*M1
      IF (LG.EQ.0) GO TO 80
      LG = -1
      JK = IT
      JP = IH2 + 1
      GO TO 80
60 LG = 1
      JK = IH2
70 IF (JK.LT.IH1) GO TO 130
      BACKSPACE 3
      BACKSPACE 3
      READ (3) E
      MF = IH1*M1
      JP = JP - IH1
80 J = JK
90 IF (J.LT.JP) GO TO 110
      MF = MF - M1
      TH = 0.
      DO 100 I=1,M1
         MQ = MF + I
         TH = TH + U(I)*E(MQ)
100 CONTINUE
      MQ = LO(J)
      U(MQ) = TH
      J = J - 1
      GO TO 90
110 IF (LG) 60, 130, 120
120 JK = JK - IH1
      GO TO 70
C FIND K, I.E. THE INDEX OF THE NEW OPTIMAL BASIC VECTOR

```

```

ARS 340
ARS 350
ARS 360
ARS 370
ARS 380
ARS 390
ARS 400
ARS 410
ARS 420
ARS 430
ARS 440
ARS 450
ARS 460
ARS 470
ARS 480
ARS 490
ARS 500
ARS 510
ARS 520
ARS 530
ARS 540
ARS 550
ARS 560
ARS 570
ARS 580
ARS 590
ARS 600
ARS 610
ARS 620
ARS 630
ARS 640
ARS 650
ARS 660
ARS 670
ARS 680
ARS 690
ARS 700
ARS 710
ARS 720
ARS 730
ARS 740
ARS 750
ARS 760
ARS 770
ARS 780
ARS 790
ARS 800
ARS 810
ARS 820
ARS 830
ARS 840
ARS 850
ARS 860
ARS 870
ARS 880
ARS 890
ARS 900
ARS 910
ARS 920
ARS 930
ARS 940
ARS 950
ARS 960
ARS 970
ARS 980
ARS 990
ARS 1000
ARS 1010
ARS 1020
ARS 1030
ARS 1040
ARS 1050
ARS 1060
ARS 1070
ARS 1080

```

```

130 CALL GEN(PROC, PRE, .FALSE., M, M1, M2, ND, K, J, RT, RT1,
* IP, NS, IBV, LO, U, D, R, RL)
C CHECK THE SOLUTION OPTIMALITY CRITERION
  IF (D(M1).LE.(-0.00001)) GO TO 230
C PRINT THE RESULTS OF ARSME
  X(M1) = -X(M1)
  WRITE (2,99999) X(M1)
  CALL ORDSNE(.TRUE., M, IBV, X)
  DO 140 I=2,M
    X(I) = X(I) + X(I-1)
140 CONTINUE
  J = 1
  K = 44
150 IF (M.LT.K) K = M
  I = 1
  JK = 1
  IF (J.GT.1) JK = K - 43
  DO 160 I=JK,K
    LO(I) = I
160 CONTINUE
  WRITE (2,99986)
  JK = J
  IF (J.GT.4) JK = 4
  GO TO (180, 190, 200, 210), JK
170 WRITE (2,99987)
  IF (K.EQ.M) RETURN
  J = J + 1
  K = K + 44
  GO TO 150
180 IF (M.GE.10) WRITE (2,99997) (LO(I),I=10,K,2)
  JK = 8
  IF (K.LT.8) JK = K
  WRITE (2,99996) (LO(I),I=1,K,2)
  WRITE (2,99995) (LO(I),I=2,JK,2)
  WRITE (2,99988)
  GO TO 220
190 IF (M.GE.46) WRITE (2,99994) (LO(I),I=46,K,2)
  WRITE (2,99993) (LO(I),I=45,K,2)
  WRITE (2,99988)
  GO TO 220
200 IF (M.GE.90) WRITE (2,99992) (LO(I),I=90,K,2)
  WRITE (2,99991) (LO(I),I=89,K,2)
  WRITE (2,99988)
  GO TO 220
210 IF (M.GE.134) WRITE (2,99990) (LO(I),I=134,K,2)
  WRITE (2,99989) (LO(I),I=133,K,2)
  WRITE (2,99988)
220 CALL GEN(PROC, .TRUE., .TRUE., M, M1, M2, ND, K, J, RT, RT1,
* IP, NS, IBV, LO, U, X, R, RL)
  GO TO 170
230 TH = 10.E70
  IF (.NOT.PRE) GO TO 250
C FIND LO(1), I.E. THE INDEX OF THE VECTOR ELIMINATED FROM THE BASIS
C IN THE FIRST ITERATION
  PRE = .FALSE.
  IT = 1
  JK = 1
  DO 240 I=1,M
    IF (D(I).EQ.0.0 .OR. TH.LE.X(I)) GO TO 240
    TH = X(I)
    L = I
240 CONTINUE
  LO(1) = L
  GO TO 380
250 IF (LG.EQ.0) GO TO 280
  LG = -1
  REWIND 3
  JP = 1 - IH1
  JK = IH1
260 IF (JK.GT.IH2) GO TO 330
270 JP = JP + IH1
  READ (3) E
280 MF = -M1
  J = JP
290 IF (J.GT.JK) GO TO 310
  MF = MF + M1

```

ARS 1090
ARS 1100
ARS 1110
ARS 1120
ARS 1130
ARS 1140
ARS 1150
ARS 1160
ARS 1170
ARS 1180
ARS 1190
ARS 1200
ARS 1210
ARS 1220
ARS 1230
ARS 1240
ARS 1250
ARS 1260
ARS 1270
ARS 1280
ARS 1290
ARS 1300
ARS 1310
ARS 1320
ARS 1330
ARS 1340
ARS 1350
ARS 1360
ARS 1370
ARS 1380
ARS 1390
ARS 1400
ARS 1410
ARS 1420
ARS 1430
ARS 1440
ARS 1450
ARS 1460
ARS 1470
ARS 1480
ARS 1490
ARS 1500
ARS 1510
ARS 1520
ARS 1530
ARS 1540
ARS 1550
ARS 1560
ARS 1570
ARS 1580
ARS 1590
ARS 1600
ARS 1610
ARS 1620
ARS 1630
ARS 1640
ARS 1650
ARS 1660
ARS 1670
ARS 1680
ARS 1690
ARS 1700
ARS 1710
ARS 1720
ARS 1730
ARS 1740
ARS 1750
ARS 1760
ARS 1770
ARS 1780
ARS 1790
ARS 1800
ARS 1810
ARS 1820
ARS 1830
ARS 1840

C CALCULATE VECTOR D	ARS 1850
L = LO(J)	ARS 1860
RR = D(L)	ARS 1870
DO 300 I=1,M	ARS 1880
MQ = MF + I	ARS 1890
D(I) = D(I) + RR*E(MQ)	ARS 1900
300 CONTINUE	ARS 1910
D(L) = D(L) - RR	ARS 1920
J = J + 1	ARS 1930
GO TO 290	ARS 1940
310 IF (LG) 320, 340, 340	ARS 1950
320 JK = JK + IH1	ARS 1960
GO TO 260	ARS 1970
330 JK = IH2 + IH	ARS 1980
LG = 1	ARS 1990
GO TO 270	ARS 2000
C INCREASE THE NUMBER OF ITERATIONS IN THE ITERATION COUNTER	ARS 2010
340 IT = IT + 1	ARS 2020
IF (IT.GT.NI) WRITE (2,99998)	ARS 2030
IF (IT.GT.NI) STOP	ARS 2040
IH = IH + 1	ARS 2050
IF (IH.LE.IH1) GO TO 360	ARS 2060
IF (LG.NE.0) GO TO 350	ARS 2070
REWIND 3	ARS 2080
WRITE (3) E	ARS 2090
LG = -1	ARS 2100
350 IH = 1	ARS 2110
IH2 = IH2 + IH1	ARS 2120
360 IF (LG.EQ.0) JK = IH	ARS 2130
C FIND LO(IT), I.E. THE INDEX OF THE VECTOR ELIMINATED FROM THE	ARS 2140
C BASIS I ITERATION IT.	ARS 2150
DO 370 I=1,M	ARS 2160
IF (D(I).LE.0.) GO TO 370	ARS 2170
RR = X(I)/D(I)	ARS 2180
IF (TH.LE.RR) GO TO 370	ARS 2190
TH = RR	ARS 2200
L = I	ARS 2210
370 CONTINUE	ARS 2220
LO(IT) = L	ARS 2230
C CALCULATE VECTORS E AND X	ARS 2240
380 MF = IH*M1 - M1	ARS 2250
MQ = MF + L	ARS 2260
RR = 1./D(L)	ARS 2270
E(MQ) = RR	ARS 2280
TH = X(L)	ARS 2290
X(L) = TH*RR	ARS 2300
DO 390 I=1,M1	ARS 2310
IF (I.EQ.L) GO TO 390	ARS 2320
MQ = MF + I	ARS 2330
G = -D(I)*RR	ARS 2340
E(MQ) = G	ARS 2350
X(I) = X(I) + G*TH	ARS 2360
390 CONTINUE	ARS 2370
IF (LG.EQ.0) GO TO 400	ARS 2380
IF (IH.GT.1) BACKSPACE 3	ARS 2390
WRITE (3) E	ARS 2400
C MEMORIZE THE INDEX OF THE VECTOR INTRODUCED INTO THE BASIS	ARS 2410
400 IBV(L) = K	ARS 2420
GO TO 40	ARS 2430
99999 FORMAT (/////35H AUTOMATIC REVISED SIMPLEX METHOD///	ARS 2440
* 17H OPTIMAL SOLUTION//33H MINIMAL SCHEDULE LENGTH TMIN =,	ARS 2450
* F12.5///)	ARS 2460
99998 FORMAT (/////37H NUMBER OF ITERATIONS GREATER THAN NI///)	ARS 2470
99997 FORMAT (30X, 1H*, 17X, 18I4)	ARS 2480
99996 FORMAT (30X, 2H* , I2, 21I4)	ARS 2490
99995 FORMAT (1H+, 31X, 4I4)	ARS 2500
99994 FORMAT (30X, 2H* , 22I4)	ARS 2510
99993 FORMAT (30X, 2H* , I2, 21I4)	ARS 2520
99992 FORMAT (30X, 2H* , 5I4, 1X, 17I4)	ARS 2530
99991 FORMAT (30X, 2H* , I2, 5I4, 1X, 16I4)	ARS 2540
99990 FORMAT (30X, 3H* , 22I4)	ARS 2550
99989 FORMAT (30X, 1H* , 22I4)	ARS 2560
99988 FORMAT (30X, 1H*/1X, 120(1H-)/30X, 1H*)	ARS 2570
99987 FORMAT (30X, 1H*/1X, 120(1H-)/)	ARS 2580
99986 FORMAT (1X, 120(1H-)/30X, 1H*/30X, 1H*, 40X, 10HACTIVITIES/	ARS 2590

IF (RT1.EQ.1) GO TO 90	GEN 720
DO 80 I=1,IC	GEN 730
KA(I) = KAS(I)	GEN 740
80 CONTINUE	GEN 750
90 CONTINUE	GEN 760
IF (DA) GO TO 120	GEN 770
IF (RT1.EQ.1) GO TO 110	GEN 780
DO 100 I=1,IC	GEN 790
KAS(I) = KA(I)	GEN 800
100 CONTINUE	GEN 810
C GENERATE SUBSETS	GEN 820
110 NR = IC	GEN 830
IF (NP.LE.(II+1)) GO TO 250	GEN 840
120 NR = II	GEN 850
130 NR = NR + 1	GEN 860
140 IF (PROC) GO TO 270	GEN 870
IF (DA) GO TO 180	GEN 880
C CHECK RESOURCE FEASIBILITY OF THE SUBSET	GEN 890
NP = 99999	GEN 900
DO 170 J=1,RT1	GEN 910
G = 0.	GEN 920
DO 150 I=1,NR	GEN 930
MQ = KAR(I)	GEN 940
MQ = KAS(MQ)	GEN 950
G = G + R(MQ,J)	GEN 960
IF (G.GT.RL(J)) GO TO 160	GEN 970
150 CONTINUE	GEN 980
GO TO 170	GEN 990
160 IF (NP.GT.I) NP = I	GEN 1000
170 CONTINUE	GEN 1010
IF (NR.GE.NP) GO TO 250	GEN 1020
C GO TO SIMPLEX PROCEDURE	GEN 1030
180 IF (PRI) GO TO 260	GEN 1040
CALL KFIND(PRE, M1, NR, K, LI, KAS, KAR, IBV, U, D)	GEN 1050
C INCREASE THE NUMBER OF SETS IN THE SET COUNTER	GEN 1060
190 LI = LI + 1	GEN 1070
IF (KAR(NR).LT.IC) GO TO 130	GEN 1080
C GENERATE THE PRIMARY SUBSET	GEN 1090
200 NR = NR - 1	GEN 1100
IF (NR.EQ.0) GO TO 300	GEN 1110
210 KAR(NR) = KAR(NR) + 1	GEN 1120
IF (KAR(NR).EQ.IC) GO TO 140	GEN 1130
MQ = NR + 1	GEN 1140
MP = NR + IC - KAR(NR)	GEN 1150
IF (PROC) GO TO 290	GEN 1160
220 DO 230 I=MQ,MP	GEN 1170
KAR(I) = KAR(I-1) + 1	GEN 1180
230 CONTINUE	GEN 1190
240 IF (KAR(NR).LE.II) NR = II - KAR(NR) + NR + 1	GEN 1200
GO TO 140	GEN 1210
C GENERATE THE NEW PRIMARY SUBSET BECAUSE OF RESOURCE INFEASIBILITY	GEN 1220
C OF THE LAST-GENERATED SET	GEN 1230
250 IF (KAR(NP).EQ.IC) GO TO 200	GEN 1240
NR = NP	GEN 1250
GO TO 210	GEN 1260
C CHECK IF THE SET COMPRISE THE OPTIMAL SOLUTION	GEN 1270
260 IF (LI.NE.IBV(JS)) GO TO 190	GEN 1280
CALL PRASET(M, NR, KW, JS, K, KAR, KAS, LO, D)	GEN 1290
JS = JS + 1	GEN 1300
IF (JS.EQ.M1) RETURN	GEN 1310
GO TO 190	GEN 1320
C GENERATE THE SETS IN THE CASE OF EQUAL RESOURCE REQUIREMENTS	GEN 1330
270 IF (NR.LE.IP) GO TO 180	GEN 1340
NR = IP	GEN 1350
IF (KAR(NR).LE.II) GO TO 280	GEN 1360
GO TO 210	GEN 1370
280 KAR(NR) = II + 1	GEN 1380
GO TO 180	GEN 1390
290 IF (MP.GT.IP) MP = IP	GEN 1400
IF (MP.LT.MQ) GO TO 240	GEN 1410
GO TO 220	GEN 1420
300 CONTINUE	GEN 1430
RETURN	GEN 1440
END	GEN 1450

```

      SUBROUTINE KFINDD(PRE, M1, NR, K, LI, KAS, KAR, IBV, U, D)
      KFI 10
C PURPOSE
      KFI 20
C SUBROUTINE KFINDD CHECKS TO DETERMINE IF THE INTRODUCTION OF THE
      KFI 30
C CURRENTLY CONSIDERED SET ( VECTOR ) INTO THE BASIS IS MORE
      KFI 40
C PROFITABLE THAN THE INTRODUCTION OF THE PREVIOUSLY GENERATED SETS.
      KFI 50
      LOGICAL PRE
      KFI 60
      DIMENSION U(M1), D(M1), IBV(M1), KAS(100), KAR(100)
      KFI 70
      IF (.NOT.(PRE .AND. (NR.EQ.1))) GO TO 10
      KFI 80
C CALCULATE THE INITIAL VALUES OF VECTOR IBV
      KFI 90
      MQ = KAR(1)
      KFI 100
      MQ = KAS(MQ)
      KFI 110
      IBV(MQ) = LI
      KFI 120
C CALCULATE VALUE G OF THE CRITERION FOR INTRODUCING A VECTOR INTO
      KFI 130
C THE BASIS
      KFI 140
      10 G = 0.
      KFI 150
      DO 20 I=1,NR
      KFI 160
      MQ = KAR(I)
      KFI 170
      MQ = KAS(MQ)
      KFI 180
      G = G + U(MQ)
      KFI 190
      20 CONTINUE
      KFI 200
      RR = 1 - NR
      KFI 210
      G = G + U(M1)*RR
      KFI 220
      IF (D(M1).LE.G) RETURN
      KFI 230
C MEMORIZE THE VECTOR FOR WHICH G IS AT MINIMUM
      KFI 240
      D(M1) = G
      KFI 250
      K = LI
      KFI 260
      MQ = M1 - 1
      KFI 270
      DO 30 I=1,MQ
      KFI 280
      D(I) = 0.
      KFI 290
      30 CONTINUE
      KFI 300
      DO 40 I=1,NR
      KFI 310
      MQ = KAR(I)
      KFI 320
      MQ = KAS(MQ)
      KFI 330
      D(MQ) = 1.
      KFI 340
      40 CONTINUE
      KFI 350
      RETURN
      KFI 360
      END
      KFI 370

```

```

      SUBROUTINE PRISET(M, NR, J, JS, K, KAR, KAS, LO, X)
      PRI 10
C PURPOSE
      PRI 20
C SUBROUTINE PRISET PRINTS THE SET OF ACTIVITIES COMPOSING THE
      PRI 30
C OPTIMAL SOLUTION AND THE TIME-INTERVALS OF ITS PERFORMANCE.
      PRI 40
      DIMENSION KAR(100), KAS(100), LO(M), X(M)
      PRI 50
      DO 10 I=1,M
      PRI 60
      LO(I) = 0
      PRI 70
      10 CONTINUE
      PRI 80
      DO 20 I=1,NR
      PRI 90
      MQ = KAR(I)
      PRI 100
      MQ = KAS(MQ)
      PRI 110
      MP = J*44
      PRI 120
      MR = MP - 44
      PRI 130
      IF (MQ.LE.MR .OR. MQ.GT.MP) GO TO 20
      PRI 140
      LO(MQ) = 1
      PRI 150
      20 CONTINUE
      PRI 160
      C = 0.
      PRI 170
      IF (JS.EQ.1) WRITE (2,99999) C, X(JS), (LO(I),I=1,K)
      PRI 180
      MQ = JS - 1
      PRI 190
      IF (JS.NE.1) WRITE (2,99999) X(MQ), X(JS), (LO(I),I=1,K)
      PRI 200
      RETURN
      PRI 210
99999 FORMAT (1X, F12.5, 3H -, F12.5, 4H * , 44I2)
      PRI 220
      END
      PRI 230

```

```

      SUBROUTINE ORDSNE(PRI, IC, KAS, RC)
      ORD 10
C PURPOSE
      ORD 20
C IF PRI IS .TRUE. SUBROUTINE ORDSNE ORDERS ELEMENTS OF VECTOR KAS
      ORD 30
C IN INCREASING ORDER OF THEIR VALUES AND VECTOR RC CORRESPONDINGLY.
      ORD 40
C IF PRI IS .FALSE. SUBROUTINE ORDSNE ORDERS ELEMENTS OF VECTOR RC
      ORD 50
C IN DECREASING ORDER OF THEIR VALUES AND VECTOR KAS CORRESPONDINGLY.
      ORD 60
      LOGICAL PRI
      ORD 70
      DIMENSION RC(IC), KAS(IC)
      ORD 80
      ID = -IC/2
      ORD 90
      10 IF (ID.GE.0) RETURN
      ORD 100
      MQ = ID + IC
      ORD 110

```

COLLECTED ALGORITHMS (cont.)**520-P12- 0**

```
DO 40 J=1,MQ
  I = J
20 IF (I.LT.1) GO TO 40
  JJ = I - ID
  IF (PRI) GO TO 50
  IF (RC(I).GE.RC(JJ)) GO TO 40
30 G = RC(I)
  RC(I) = RC(JJ)
  RC(JJ) = G
  II = KAS(I)
  KAS(I) = KAS(JJ)
  KAS(JJ) = II
  I = I + ID
  GO TO 20
40 CONTINUE
  ID = ID/2
  GO TO 10
50 IF (KAS(I).LE.KAS(JJ)) GO TO 40
  GO TO 30
END
```

ORD	120
ORD	130
ORD	140
ORD	150
ORD	160
ORD	170
ORD	180
ORD	190
ORD	200
ORD	210
ORD	220
ORD	230
ORD	240
ORD	250
ORD	260
ORD	270
ORD	280
ORD	290
ORD	300
ORD	310

ALGORITHM 521

Repeated Integrals of the Coerror Function [S15]

WALTER GAUTSCHI
Purdue University

Key Words and Phrases: repeated integrals of the coerror function, Taylor's series, three-term recurrence relation, Miller's backward recurrence algorithm
CR Categories: 5.12
Language: Fortran

DESCRIPTION

This algorithm implements the procedures developed in [1].

REFERENCES

- GAUTSCHI, W. Evaluation of the repeated integrals of the coerror function. *ACM Trans. Math. Software* 3, 3 (Sept. 1977), 240-252.

ALGORITHM

```

      SUBROUTINE INERFC(X, NMAX, ACC, FZERO, F, IFLAG)           INE  10
C THIS PROGRAM GENERATES FN(X) FOR N=1,2,...,NMAX, WHERE      INE  20
C FN(X)=EXP(X**2)*INERFC(X), IF X.GE.0., AND FN(X)=INERFC(X), IF INE  30
C X.LT.0., INERFC(X) BEING THE N-TH REPEATED INTEGRAL OF THE COERROR INE  40
C FUNCTION ERFC(X), EXTENDED FROM X TO INFINITY. THE PROGRAM ALSO INE  50
C SUPPLIES EXP(X**2)*ERFC(X), IF X.GE.0., AND ERFC(X), IF X.LT.0.. INE  60
C THE USER HAS TO SPECIFY THE DESIRED ACCURACY, AS WELL AS THE PRECISION INE  70
C OF THE HIGHEST PRECISION MODE HE IS PREPARED TO USE. THE FORMER IS INE  80
C INPUT VIA THE PARAMETER ACC IN THE CALL LIST, WHICH DESIGNATES THE INE  90
C NUMBER OF CORRECT SIGNIFICANT DECIMAL DIGITS DESIRED IN THE RESULTS. INE 100
C THE LATTER IS INPUT VIA THE PARAMETER PREC IN THE DATA STATEMENT. THE INE 110
C VALUE OF PREC SHOULD BE SET APPROXIMATELY EQUAL TO BETA*ALOG(2.)/ INE 120
C ALOG(10.), WHERE BETA IS THE NUMBER OF BINARY DIGITS AVAILABLE IN THE INE 130
C MANTISSA OF THE HIGHEST PRECISION FLOATING-POINT WORD. ORDINARILY, THE INE 140
C HIGHEST PRECISION IS DOUBLE PRECISION, EITHER HARDWARE GENERATED OR INE 150
C SOFTWARE GENERATED, BUT IT MAY ALSO BE SINGLE PRECISION. IN ANY CASE, INE 160
C ACC HAS TO BE LESS THAN PREC, THE DIFFERENCE BEING AT LEAST EQUAL TO INE 170
C ONE, PREFERABLY SEVERAL UNITS LARGER. EVIDENTLY, ACC SHOULD NOT BE INE 180
C SPECIFIED TO EXCEED THE PRECISION OF THE LOWEST PRECISION MODE USED. INE 190
C THE DATA STATEMENT CONTAINS ANOTHER MACHINE-DEPENDENT PARAMETER, INE 200
C BOTEXP, WHICH IS, APPROXIMATELY, THE SMALLEST NEGATIVE NUMBER SUCH INE 210
C THAT 10.**BOTEXP IS STILL REPRESENTABLE ON THE COMPUTER IN SINGLE PRE- INE 220
C CISION FLOATING-POINT ARITHMETIC. THE TYPE OF BOTEXP IS REAL. IN THE INE 230
C DATA STATEMENT BELOW THE PARAMETERS PREC AND BOTEXP ARE SET TO CORRE- INE 240
C SPOND TO THE MACHINE CHARACTERISTICS OF THE CDC 6500 COMPUTER. INE 250
C PARAMETER LIST INE 260
C X -- THE ARGUMENT OF INERFC. TYPE REAL. INE 270
C NMAX -- THE UPPER LIMIT ON N. TYPE INTEGER. INE 280
C ACC -- THE NUMBER OF CORRECT SIGNIFICANT DECIMAL DIGITS DESIRED INE 290
C IN THE RESULTS. TYPE REAL. INE 300

```

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```

C FZERO -- AN OUTPUT VARIABLE RETURNING EXP(X**2)*ERFC(X), IF INE 310
C X.GE.0., AND ERFC(X), IF X.LT.0. TYPE REAL. INE 320
C F -- THE NAME OF A ONE-DIMENSIONAL ARRAY HOLDING THE RESULTS INE 330
C FN(X), N=1,2,...,NMAX. THE ARRAY HAS DIMENSION NMAX. INE 340
C IFLAG -- AN ERROR FLAG INDICATING A NUMBER OF FATAL ERROR CON- INE 350
C DITIONS UPON EXECUTION. TYPE INTEGER. THE VALUES OF THIS INE 360
C VARIABLE HAVE THE FOLLOWING MEANINGS. INE 370
C IFLAG=0 -- NO ERROR CONDITION. INE 380
C IFLAG=1 -- ILLEGAL NEGATIVE OR ZERO VALUE OF NMAX. INE 390
C IFLAG=2 -- EXCESSIVELY LARGE VALUE OF NMAX. REDIMENSION THE INE 400
C ARRAYS R0 AND R1 TO HAVE DIMENSION .GE. NMAX, IF INE 410
C INDEED NMAX MUST BE THAT LARGE. INE 420
C IFLAG=3 -- INAPPROPRIATE ACCURACY SPECIFICATION. INE 430
C IFLAG=4 -- UNUSUAL UNDERFLOW OF EXP(-X**2). TRY REDUCING INE 440
C THE ERROR TOLERANCE PREC-ACC, EITHER BY DE- INE 450
C CREATING PREC OR BY INCREASING ACC. INE 460
C IFLAG=5 -- TAYLOR SERIES FOR ERFC(X) DOES NOT CONVERGE INE 470
C WITHIN 200 TERMS FOR SOME UNKNOWN REASON. INE 480
C IFLAG=6 -- NU IN THE BACKWARD RECURRENCE ALGORITHM EXCEEDS INE 490
C 5000, PROBABLY BECAUSE X IS A SMALL POSITIVE INE 500
C NUMBER, NMAX RELATIVELY LARGE, AND THE ERROR INE 510
C TOLERANCE SMALL. TRY INCREASING THE ERROR TOLER- INE 520
C ANCE PREC-ACC, EITHER BY INCREASING PREC OR BY INE 530
C DECREASING ACC. INE 540
C THE SUBROUTINE BELOW IS WRITTEN WITH THE ASSUMPTION THAT PART OF INE 550
C THE COMPUTATION IS DONE IN DOUBLE PRECISION AND THE REST IN SINGLE INE 560
C PRECISION. TO OBTAIN A VERSION OF THIS SUBROUTINE, THAT OPERATES INE 570
C ENTIRELY IN SINGLE PRECISION, THE FOLLOWING CHANGES MUST BE MADE. INE 580
C (1) DELETE THE DOUBLE PRECISION TYPE DECLARATION. INE 590
C (2) REPLACE THE SECTION OF PROGRAM FROM STATEMENT LABELED 30 INE 600
C (INCLUSIVE) TO STATEMENT LABELED 60 (INCLUSIVE) BY THE FOLLOWING INE 610
C PIECE OF PROGRAM. INE 620
C 30 DEFS=.5*10.**(-PREC) INE 630
C TE=C*X INE 640
C SUM=1.-TE INE 650
C N=0 INE 660
C N1=1 INE 670
C 40 N=N+1 INE 680
C N1=N1+2 INE 690
C IF(N.GT.200) GOTO 270 INE 700
C TE=-TE*XSQ/FLOAT(N) INE 710
C TERM=TE/FLOAT(N1) INE 720
C SUM=SUM-TERM INE 730
C IF(ABS(TERM).GT.DEFS*ABS(SUM)) GOTO 40 INE 740
C F1=T*SUM INE 750
C THE ARRAYS R0 AND R1 IN THE DIMENSION STATEMENT ARE LOCAL TO THE INE 760
C SUBROUTINE. THEY BOTH NEED TO HAVE DIMENSION LARGER OR EQUAL TO NMAX. INE 770
C WE DECLARED THEIR DIMENSION TO BE 500, ASSUMING THAT NMAX WILL NOT INE 780
C EXCEED 500. IN THE EVENT IT DOES, THE SUBROUTINE EXITS WITH THE ERROR INE 790
C FLAG AT 2. IF THE USER EXPECTS HIS VALUES OF NMAX TO BE CONSISTENTLY INE 800
C MUCH LESS THAN 500, HE CAN SAVE STORAGE BY LOWERING THE DIMENSION OF INE 810
C R0 AND R1 AND AT THE SAME TIME ADJUSTING THE ERROR EXIT STATEMENT (THE INE 820
C THIRD EXECUTABLE STATEMENT IN THE PROGRAM). INE 830
C WITH THE EXCEPTION NOTED UNDER IFLAG=4, ANY VARIABLE THAT UNDER- INE 840
C FLOWS MAY BE SET EQUAL TO ZERO. NO PRECAUTION IS TAKEN AGAINST OVER- INE 850
C FLOW, WHICH MAY OCCUR IN THE DO-LOOP AFTER STATEMENT 70 IF X.LT.0. INE 860
C AND ABS(X) AND/OR NMAX IS LARGE. INE 870
C REFERENCE - W. GAUTSCHI, EVALUATION OF THE REPEATED INTEGRALS OF INE 880
C THE COERROR FUNCTION, ACM TRANS. MATH. SOFTWARE. INE 890
C DIMENSION F(NMAX), R0(500), R1(500) INE 900
C DOUBLE PRECISION DEFS, DC, DX, DXSQ, DFM1, DF0, DF1, DN, DN1, INE 910
C * DTE, DTERM, DSUM INE 920
C LOGICAL FRSTTM INE 930
C DATA FRSTTM, PREC, BOTEXP /.TRUE.,28.8989,-293./ INE 940
C IFLAG = 0 INE 950
C IF (NMAX.LT.1) GO TO 230 INE 960
C IF (NMAX.GT.500) GO TO 240 INE 970
C TOL = PREC - ACC INE 980
C IF (TOL.LT.1.) GO TO 250 INE 990
C I = 1 INE 1000
C EPS = .5*10.**(-ACC) INE 1010
C XSQ = X*X INE 1020
C IF (.NOT.FRSTTM) GO TO 10 INE 1030
C P = ATAN(1.) INE 1040
C AL10 = ALOG(10.) INE 1050
C C = SQRT(1./P) INE 1060
C DC = DSQRT(1.D0/DATAN(1.D0)) INE 1070

```

```

FRSTTM = .FALSE.
10 S = 0.
IF (-XSQ.GT.AL10*BOTEXP) S = EXP(-XSQ)
B = .5*AL10*TOL + .25*ALOG(2.*P)
B1 = .5*AL10*(TOL-1.)
IF (XSQ.GT.B) GO TO 90
IF (X.LT.0.) GO TO 20
F0 = C
IF (S.EQ.0.) GO TO 260
T = 1./S
SQ = SQRT(2.*FLOAT(NMAX))
IF ((X.GE.1.12 .AND. XSQ+SQ*X.GT.B) .OR. (X.LT.1.12 .AND.
* SQ*X.GT.B1)) GO TO 100
GO TO 30
20 F0 = C*S
T = 1.
C
C EVALUATION OF ERFC(X) BY TAYLOR SERIES IN DOUBLE PRECISION
C
30 DEPS = .5D0*10.D0**(-PREC)
DX = DBLE(X)
DXSQ = DX*DX
DTE = DC*DX
DSUM = 1.D0 - DTE
DN = 0.D0
DN1 = 1.D0
40 DN = DN + 1.D0
DN1 = DN1 + 2.D0
IF (DN.GT.2.D2) GO TO 270
DTE = -DTE*DXSQ/DN
DTERM = DTE/DN1
DSUM = DSUM - DTERM
IF (DABS(DTERM).GT.DEPS*DABS(DSUM)) GO TO 40
IF (X.LT.0.) GO TO 60
C
C EVALUATION OF EXP(X**2)*INERFC(X),X.GE.0.,BY FORWARD RECURSION IN
C DOUBLE PRECISION
C
DF0 = DC
DF1 = DEXP(DXSQ)*DSUM
FZERO = SNGL(DF1)
DO 50 N=1,NMAX
DFM1 = DF0
DF0 = DF1
DF1 = (-DX*DF0+.5D0*DFM1)/DBLE(FLOAT(N))
F(N) = SNGL(DF1)
50 CONTINUE
RETURN
C
C EVALUATION OF INERFC(X),X.LT.0.,BY FORWARD RECURSION
C
60 F1 = SNGL(DSUM)
70 FZERO = F1
DO 80 N=1,NMAX
FM1 = F0
F0 = F1
F1 = (-X*F0+.5*FM1)/FLOAT(N)
F(N) = F1
80 CONTINUE
RETURN
90 IF (X.LT.0.) GO TO 110
100 N0 = NMAX
X0 = X
GO TO 130
110 IF (XSQ.LI.(ACC+1.)*AL10-.572) GO TO 120
F0 = C*S
F1 = 2.
GO TO 70
120 I = 2
N0 = 1
X0 = -X
C
C BACKWARD RECURRENCE ALGORITHM FOR EVALUATING FN(ABS(X)) FOR N=1,2,
C ...,N0, WHERE N0=NMAX IF X.GT.0. AND N0=1 IF X.LT.0.
C
130 DO 140 N=1,N0
R1(N) = 0.

```

INE 1080
INE 1090
INE 1100
INE 1110
INE 1120
INE 1130
INE 1140
INE 1150
INE 1160
INE 1170
INE 1180
INE 1190
INE 1200
INE 1210
INE 1220
INE 1230
INE 1240
INE 1250
INE 1260
INE 1270
INE 1280
INE 1290
INE 1300
INE 1310
INE 1320
INE 1330
INE 1340
INE 1350
INE 1360
INE 1370
INE 1380
INE 1390
INE 1400
INE 1410
INE 1420
INE 1430
INE 1440
INE 1450
INE 1460
INE 1470
INE 1480
INE 1490
INE 1500
INE 1510
INE 1520
INE 1530
INE 1540
INE 1550
INE 1560
INE 1570
INE 1580
INE 1590
INE 1600
INE 1610
INE 1620
INE 1630
INE 1640
INE 1650
INE 1660
INE 1670
INE 1680
INE 1690
INE 1700
INE 1710
INE 1720
INE 1730
INE 1740
INE 1750
INE 1760
INE 1770
INE 1780
INE 1790
INE 1800
INE 1810
INE 1820
INE 1830
INE 1840

140	CONTINUE	INE 1850
	FN0 = FLOAT(N0)	INE 1860
	NU = IFIX((SQRT(FN0)+(2.3026*ACC+1.3863)/(2.8284*X0))**2-5.)	INE 1870
150	NU = NU + 10	INE 1880
	IF (NU.GT.5000) GO TO 280	INE 1890
	N0P1 = N0 + 1	INE 1900
	NUM = NU - N0P1	INE 1910
	DO 160 N=1,N0	INE 1920
	R0(N) = R1(N)	INE 1930
160	CONTINUE	INE 1940
	R = 0.	INE 1950
	DO 170 K=1,NUM	INE 1960
	N = NU - K	INE 1970
	R = .5/(X0+FLOAT(N+1)*R)	INE 1980
170	CONTINUE	INE 1990
	DO 180 K=1,N0	INE 2000
	N = N0P1 - K	INE 2010
	R = .5/(X0+FLOAT(N+1)*R)	INE 2020
	R1(N) = R	INE 2030
180	CONTINUE	INE 2040
	DO 190 N=1,N0	INE 2050
	IF (ABS(R1(N)-R0(N)).GT.EPS*ABS(R1(N))) GO TO 150	INE 2060
190	CONTINUE	INE 2070
	F0 = .5*C/(X0+R1(1))	INE 2080
	FZERO = F0	INE 2090
	FF = F0	INE 2100
	DO 200 N=1,N0	INE 2110
	FF = R1(N)*FF	INE 2120
	F(N) = FF	INE 2130
200	CONTINUE	INE 2140
	GO TO (210, 220), I	INE 2150
210	RETURN	INE 2160
C		INE 2170
C	STARTING VALUE FOR INERFC(X),X.LT.0.,PRIOR TO FORWARD RECURSION	INE 2180
C		INE 2190
220	F1 = 2. - S*F0	INE 2200
	F0 = C*S	INE 2210
	GO TO 70	INE 2220
230	IFLAG = 1	INE 2230
	RETURN	INE 2240
240	IFLAG = 2	INE 2250
	RETURN	INE 2260
250	IFLAG = 3	INE 2270
	RETURN	INE 2280
260	IFLAG = 4	INE 2290
	RETURN	INE 2300
270	IFLAG = 5	INE 2310
	RETURN	INE 2320
280	IFLAG = 6	INE 2330
	RETURN	INE 2340
	END	INE 2350

ALGORITHM 522

ESOLVE, Congruence Techniques for the Exact Solution of Integer Systems of Linear Equations [F4]

S. CABAY and T. P. L. LAM
University of Alberta

Key Words and Phrases: symbolic and algebraic manipulation, systems of linear equations, congruence techniques, modular methods, Chinese remainder theorem
CR Categories: 5.14
Language: Fortran

DESCRIPTION

Descriptions of the main algorithm, ESOLVE, and accompanying subroutines SUBBND, MRADIX, and FRADIX, together with experimental results, are given in [1].

REFERENCES

1. CABAY, S., AND LAM, T.P.L. Congruence techniques for the exact solution of integer systems of linear equations. *ACM Trans. Math. Software* 3, 4 (Dec. 1977), 386-497.

ALGORITHM

```

C
C*****
C
C          TWO TEST EXAMPLES
C          -----
C
C*****
C
C          INTEGER P(100),AMOD(20,21),Y(32,20,1),DET(32),T,TWOT1,
1          SUM(7),GAMMA(7),DELTA(7),A(10),C(10),B
C          INTEGER A1(1,5,5),B1(1,5,1),A2(3,20,20),B2(3,20,1),TEMP1,Q
C          REAL S(21)
C          COMMON /PRIMEB/ P, IPRIME(100)
C          B = 10000
C
C*****
C
C          EXAMPLE 1 - PASCAL MATRIX.
C
C
C          MAIN0010
C          MAIN0020
C          MAIN0030 *
C          MAIN0040 *
C          MAIN0050 *
C          MAIN0060 *
C          MAIN0070
C          MAIN0080
C          MAIN0090
C          MAIN0100
C          MAIN0110
C          MAIN0120
C          MAIN 130
C          MAIN0140
C          MAIN0150
C          MAIN0160
C          MAIN0170 *
C          MAIN0180 *
C          MAIN0190 *

```

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```

C           A1*X = B1.           *           MAIN0200
C                                           *           MAIN0210
C           T = 1                 *           MAIN0220
C           N = 5                 *           MAIN0230
C           M = 1                 *           MAIN0240
C                                           *           MAIN0250
C*****                               *           MAIN0260
C                                           *           MAIN0270
C           T = 1                 *           MAIN0280
C           N = 5                 *           MAIN0290
C           M = 1                 *           MAIN0300
C           NM = N+M              *           MAIN0310
C                                           *           MAIN0320
C           **GENERATE THE MATRICES A1 AND B1. *           MAIN0330
C                                           *           MAIN0340
C           DO 1 I=1,N             *           MAIN0350
C             A1(1,I,1) = 1        *           MAIN0360
C           1     A1(1,1,I) = 1    *           MAIN0370
C           C           OD         *           MAIN0380
C           DO 2 I=2,N             *           MAIN0390
C             DO 2 J=2,N           *           MAIN0400
C           2     A1(1,I,J) = A1(1,I,J-1)+A1(1,I-1,J) *           MAIN0410
C           C           OD         *           MAIN0420
C           OD                     *           MAIN0430
C           B1(1,1,1) = 1         *           MAIN0440
C           DO 3 I=2,N            *           MAIN0450
C           3     B1(1,I,1) = B1(1,I-1,1)+A1(1,I,N) *           MAIN0460
C           C           OD         *           MAIN0470
C                                           *           MAIN0480
C           **PRINT A1 AND B1 IN FIXED-RADIX FORM. *           MAIN0490
C                                           *           MAIN0500
C           WRITE(6,42)           *           MAIN0510
C           DO 5 K=1,T             *           MAIN0520
C             WRITE(6,46) K       *           MAIN0530
C             DO 4 I=1,N          *           MAIN0540
C           4     WRITE(6,43) (A1(K,I,J),J=1,N) *           MAIN0550
C           C           OD         *           MAIN0560
C                                           *           CONTINUEMAIN0570
C           OD                     *           MAIN0580
C           WRITE(6,44)           *           MAIN0590
C           DO 7 K=1,T             *           MAIN0600
C             WRITE(6,49) K       *           MAIN0610
C             DO 6 I=1,N          *           MAIN0620
C           6     WRITE(6,43) (B1(K,I,J),J=1,M) *           MAIN0630
C           C           OD         *           MAIN0640
C                                           *           CONTINUEMAIN0650
C           OD                     *           MAIN0660
C                                           *           MAIN0670
C           **CONVERT A1 AND B1 FROM FIXED-RADIX TO MIXED-RADIX FORM. *           MAIN0680
C                                           *           MAIN0690
C           N1 = 1                 *           MAIN0700
C           TEMP = (FLOAT(N1)*ALOG(FLOAT(B)) + ALOG(2.0))/ALOG(FLOAT(P(1))) *           MAIN0710
C           LMAX = TEMP            *           MAIN0720
C           IF (TEMP-FLOAT(LMAX).GT.0.01) LMAX = LMAX+1 *           MAIN0730
C           DO 10 I=1,N            *           MAIN0740
C             DO 10 J=1,N          *           MAIN0750
C             DO 8 K=1,N1         *           MAIN0760
C           8     A(K) = A1(K,I,J) *           MAIN0770
C           C           OD         *           MAIN0780
C             CALL MRADIX(A,N1,C,LMAX,L,B,IER) *           MAIN0790
C             DO 9 K=1,L          *           MAIN0800
C           9     A1(K,I,J) = C(K) *           MAIN0810
C           C           OD         *           MAIN0820
C           OD                     *           MAIN0830
C           OD                     *           MAIN0840
C                                           *           CONTINUEMAIN0850
C           DO 13 I=1,N            *           MAIN0860
C             DO 13 J=1,M          *           MAIN0870
C             DO 11 K=1,N1        *           MAIN0880
C           11     A(K) = B1(K,I,J) *           MAIN0890
C           C           OD         *           MAIN0900
C             CALL MRADIX(A,N1,C,LMAX,L,B,IER) *           MAIN0910
C             DO 12 K=1,L         *           MAIN0920
C           12     B1(K,I,J) = C(K) *           MAIN0930
C           C           OD         *           MAIN0940

```

```

C          OD                                MAIN0950
C          OD                                MAIN0960
13        CONTINUEMAIN0970
C          MAIN0980
C          **PRINT A1 AND B1 IN MIXED-RADIX FORM.  MAIN0990
C          MAIN1000
C          WRITE(6,45)                        MAIN1010
C          DO 15 K=1,T                          MAIN1020
C            WRITE(6,46) K                      MAIN1030
C            DO 14 I=1,N                        MAIN1040
14          WRITE(6,47) (A1(K,I,J),J=1,N)      MAIN1050
C          OD                                  MAIN1060
15        CONTINUEMAIN1070
C          OD                                  MAIN1080
C          WRITE(6,48)                          MAIN1090
C          DO 17 K=1,T                          MAIN1100
C            WRITE(6,49) K                      MAIN1110
C            DO 16 I=1,N                        MAIN1120
16          WRITE(6,47) (B1(K,I,J),J=1,M)      MAIN1130
C          OD                                  MAIN1140
17        CONTINUEMAIN1150
C          OD                                  MAIN1160
C          MAIN1170
C          **COMPUTE MAXPRM.                    MAIN1180
C          MAIN1190
C          TWOT1 = 2*T+1                         MAIN1200
C          CALL SUBBND(A1,B1,GAMMA,DELTA,S,SUM,T,N,M,NM, MAIN1210
1          TWOT1,BOUND,MAXPRM)                 MAIN1220
C          WRITE(6,50) BOUND,MAXPRM            MAIN1230
C          MAIN1240
C          **SOLVE THE SYSTEM OF LINEAR EQUATIONS A1*X=B1. MAIN1250
C          MAIN1260
C          CALL DRIVER(A1,B1,AMOD,T,N,M,NM,Y,DET,MAXPRM,B) MAIN1270
C          MAIN1280
C          *****                             MAIN1290
C          *                                     *
C          EXAMPLE 2 - PASCAL MATRIX.           *
C          *                                     *
C          A2*X = B2.                           *
C          *                                     *
C          T = 3                                 *
C          N = 20                                *
C          M = 1                                 *
C          *                                     *
C          *****                             MAIN1390
C          T = 3                                 MAIN1400
C          N = 20                                MAIN1410
C          M = 1                                 MAIN1420
C          NM = N+M                             MAIN1430
C          MAIN1440
C          MAIN1450
C          **GENERATE THE MATRICES A2 AND B2.    MAIN1460
C          MAIN1470
C          DO 18 I=1,N                           MAIN1480
C            A2(3,I,1) = 1                       MAIN1490
18          A2(3,1,I) = 1                       MAIN1500
C          OD                                    MAIN1510
C          DO 19 K=1,2                           MAIN1520
C            DO 19 I=1,N                         MAIN1530
C              A2(K,I,1) = 0                    MAIN1540
19          A2(K,1,I) = 0                      MAIN1550
C          OD                                    MAIN1560
C          OD                                    MAIN1570
C          DO 21 I=2,N                           MAIN1580
C            DO 21 J=2,N                         MAIN1590
C              Q = 0                            MAIN1600
C              TEMP1 = 0                        MAIN1610
C              DO 20 K=1,T                      MAIN1620
C                K1 = T-K+1                    MAIN1630
C                TEMP1 = A2(K1,I,J-1)+A2(K1,I-1,J)+Q MAIN1640
C                Q = TEMP1/B                   MAIN1650
20          A2(K1,I,J) = TEMP1 - Q*B           MAIN1660
C          OD                                    MAIN1670
C          OD                                    MAIN1680
C          OD                                    MAIN1690

```

```

21          B2(1,1,1) = 0
           B2(2,1,1) = 0
           B2(3,1,1) = 1
           DO 23 I=2,N
             Q = 0
             TEMP1 = 0
             DO 22 K=1,T
               K1 = T-K+1
               TEMP1 = B2(K1,I-1,1)+A2(K1,I,N)+Q
               Q = TEMP1/B
           B2(K1,I,1) = TEMP1 - Q*B
22
C           OD
C           OD
23
C           **PRINT A2 AND B2 IN FIXED-RADIX FORM.
C
           WRITE(6,42)
           DO 25 K=1,T
             WRITE(6,46) K
             K1 = T-K+1
           DO 24 I=1,N
             WRITE(6,43) (A2(K1,I,J),J=1,N)
24
C           OD
C           OD
25
           WRITE(6,44)
           DO 27 K=1,T
             WRITE(6,49) K
             K1 = T-K+1
           DO 26 I=1,N
             WRITE(6,43) (B2(K1,I,J),J=1,M)
26
C           OD
C           OD
27
C           **CONVERT A2 AND B2 FROM FIXED-RADIX TO MIXED-RADIX FORM.
C
           N1 = 3
           TEMP = (FLOAT(N1)*ALOG(FLOAT(B)) + ALOG(2.0))/ALOG(FLOAT(P(1)))
           LMAX = TEMP
           IF (TEMP-FLOAT(LMAX).GT.0.01) LMAX = LMAX+1
           DO 32 I=1,N
             DO 32 J=1,N
               DO 28 K=1,N1
                 A(K) = A2(K,I,J)
28
C           OD
           CALL MRADIX(A,N1,C,LMAX,L,B,IER)
           DO 29 K=1,L
             A2(K,I,J) = C(K)
29
C           OD
           IF (L.LT.N1)
C             THEN
30             L1 = L+1
             DO 31 K=L1,N1
               A2(K,I,J) = 0
31
C           OD
C           OD
C           OD
32
           DO 37 I=1,N
             DO 37 J=1,M
               DO 33 K=1,N1
                 A(K) = B2(K,I,J)
33
C           OD
           CALL MRADIX(A,N1,C,LMAX,L,B,IER)
           DO 34 K=1,L
             B2(K,I,J) = C(K)
34
C           OD
           IF (L.LT.N1)
C             THEN
35             L1 = L+1

```

```

CONTINUEMAIN1700
MAIN1710
MAIN1720
MAIN1730
MAIN1740
MAIN1750
MAIN1760
MAIN1770
MAIN1780
MAIN1790
MAIN1800
MAIN1810
MAIN1820
MAIN1830
CONTINUEMAIN1840
MAIN1850
MAIN1860
MAIN1870
MAIN1880
MAIN1890
MAIN1900
MAIN1910
MAIN1920
MAIN1930
MAIN1940
MAIN1950
CONTINUEMAIN1960
MAIN1970
MAIN1980
MAIN1990
MAIN2000
MAIN2010
MAIN2020
MAIN2030
MAIN2040
CONTINUEMAIN2050
MAIN2060
MAIN2070
MAIN2080
MAIN2090
MAIN2100
MAIN2110
MAIN2120
MAIN2130
MAIN2140
MAIN2150
MAIN2160
MAIN2170
MAIN2180
MAIN2190
MAIN2200
MAIN2210
GOTO 30 MAIN2220
GOTO 32 MAIN2230
MAIN2240
MAIN2250
MAIN2260
MAIN2270
MAIN2280
MAIN2290
MAIN2300
CONTINUEMAIN2310
MAIN2320
MAIN2330
MAIN2340
MAIN2350
MAIN2360
MAIN2370
MAIN2380
MAIN2390
MAIN2400
GOTO 35 MAIN2410
GOTO 37 MAIN2420
MAIN2430
MAIN2440

```

```

DO 36 K=L1,N1                                MAIN2450
  B2(K,I,J) = 0                                MAIN2460
C                                               MAIN2470
C       OD                                    MAIN2480
C     OD                                     MAIN2490
37 CONTINUEMAIN2500
C                                               MAIN2510
C     **PRINT A2 AND B2 IN MIXED-RADIX FORM.   MAIN2520
C                                               MAIN2530
C     WRITE(6,45)                             MAIN2540
C     DO 39 K=1,T                              MAIN2550
C     WRITE(6,46) K                            MAIN2560
C     DO 38 I=1,N                              MAIN2570
38 WRITE(6,47) (A2(K,I,J),J=1,N)             MAIN2580
C     OD                                       MAIN2590
39 CONTINUEMAIN2600
C     OD                                       MAIN2610
C     WRITE(6,48)                             MAIN2620
C     DO 41 K=1,T                              MAIN2630
C     WRITE(6,49) K                            MAIN2640
C     DO 40 I=1,N                              MAIN2650
40 WRITE(6,47) (B2(K,I,J),J=1,M)           MAIN2660
C     OD                                       MAIN2670
41 CONTINUEMAIN2680
C     OD                                       MAIN2690
C                                               MAIN2700
C     **COMPUTE MAXPRM.                       MAIN2710
C                                               MAIN2720
C     TWOT1 = 2*T+1                            MAIN2730
C     CALL SUBBND(A2,B2,GAMMA,DELTA,S,SUM,T,N,M,NM,
1     TWOT1,BOUND,MAXPRM)                   MAIN2740
C     WRITE(6,50) BOUND,MAXPRM                MAIN2750
C                                               MAIN2760
C                                               MAIN2770
C     **SOLVE THE SYSTEM OF LINEAR EQUATIONS A2*X=B2. MAIN2780
C                                               MAIN2790
C     CALL DRIVER(A2,B2,AMOD,T,N,M,NM,Y,DET,MAXPRM,B) MAIN2800
C                                               MAIN2810
42 FORMAT(1H1,16X,5HINPUT/17X,5(1H*)/5X,    MAIN2820
1     38HFIXED-RADIX REPRESENTATION, BASE=10000/10X, MAIN2830
2     40HA = A(1,N,N)+A(2,N,N)*BASE+...+A(T,N,N)*, MAIN2840
3     11HBASE**(T-1))                       MAIN2850
43 FORMAT(5X,20(I4,2X))                     MAIN2860
44 FORMAT(/10X,40HB = B(1,N,M)+B(2,N,M)*BASE+...+B(T,N,M)*,
1     11HBASE**(T-1))                       MAIN2880
45 FORMAT(/5X,26HMIXED-RADIX REPRESENTATION/10X,
1     40HA = A(1,N,N)+A(2,N,N)*P(1)+...+A(T,N,N)*, MAIN2890
2     15HP(1)*...*P(T-1))                 MAIN2910
46 FORMAT(/10X,2HA(I1,5H,N,N))             MAIN2920
47 FORMAT(1H ,20(I6))                       MAIN2930
48 FORMAT(/10X,40HB = B(1,N,M)+B(2,N,M)*P(1)+...+B(T,N,M)*,
1     15HP(1)*...*P(T-1))                 MAIN2940
49 FORMAT(/10X,2HB(I1,5H,N,M))             MAIN2950
50 FORMAT(/10X,13HLOG(BOUND) = ,F8.4,2X,9HMAXPRM = ,I3)
STOP                                         MAIN2960
END                                          MAIN2970
                                           MAIN2980
                                           MAIN2990

C
C*****
C
C     CALL ESOLVE.                            *
C     PRINT OUTPUTS IN MIXED-RADIX FORM.     *
C     CALL FRADIX.                            *
C     PRINT OUTPUTS IN FIXED-RADIX FORM.     *
C*****
C
C     SUBROUTINE DRIVER(A,B,AMOD,T,N,M,NM,Y,DET,MAXPRM,BASE)
C     INTEGER T,A(T,N,N),B(T,N,M),AMOD(N,NM),Y(MAXPRM,N,M),
1     DET(MAXPRM),BASE,P(100)
C     INTEGER A1(10),C1(10)
C     COMMON /PRIMEB/ P, IPRIME(100)
C     LOGICAL RTEST
C
C     **SOLVE THE SYSTEM OF LINEAR EQUATIONS AX=B.

```

```

DRIV0010
DRIV0020
DRIV0030
DRIV0040
DRIV0050
DRIV0060
DRIV0070
DRIV0080
DRIV0090
DRIV0100
DRIV0110
DRIV0120
DRIV0130
DRIV 460
DRIV 140
DRIV0150
DRIV0160
DRIV0170

```

```

C          RTEST = .TRUE.
          CALL ESOLVE(A,B,AMOD,T,N,M,NM,DET,Y,MAXPRM,NOCOEF,
1          RTEST,IER)
C          **PRINT OUTPUT PARAMETERS.
C          WRITE(6,6) IER,NOCOEF
C          **PRINT Y IN MIXED-RADIX FORM.
C          WRITE(6,7)
          DO 2 J=1,M
            WRITE(6,8) J
            DO 1 I=1,N
1              WRITE(6,10) (Y(K,I,J),K=1,NOCOEF)
C              OD
2          OD
C          **PRINT DET IN MIXED-RADIX FORM.
C          WRITE(6,9)
          WRITE(6,10) (DET(K),K=1,NOCOEF)
C          **CONVERT Y AND DET FROM MIXED-RADIX TO FIXED-RADIX FORM.
C          **PRINT Y AND DET IN FIXED-RADIX FORM.
C          TEMP = (FLOAT(NOCOEF)*ALOG(FLOAT(P(NOCOEF))) - ALOG(2.0)) /
1          ALOG(FLOAT(BASE))
          LMAX = TEMP
          IF (TEMP-FLOAT(LMAX).GT.0.01) LMAX = LMAX+1
          WRITE(6,11)
          DO 5 J=1,M
            WRITE(6,8) J
            DO 4 I=1,N
3              DO 3 K=1,NOCOEF
C                  C1(K) = Y(K,I,J)
C              OD
4              CALL FRADIX(C1,NOCOEF,A1,LMAX,L,BASE,IER)
          WRITE(6,10) (A1(II),II=1,L)
C          OD
5          CONTINUEDRIV0610
C          OD
          CALL FRADIX(DET,NOCOEF,A1,LMAX,L,BASE,IER)
          WRITE(6,12) (A1(I),I=1,L)
C
6          FORMAT(/16X,7HOUTPUTS/16X,7(1H*)/6X,
1          6HIER = ,I3,5X,9HNOCOEF = ,I3)
7          FORMAT(/16X,21HY IN MIXED-RADIX FORM/10X,
1          35HFROM LEFT TO RIGHT, THE NUMBERS ARE,
2          25H COEF(1),...,COEF(L), AND/10X,
3          33HY(I,J) = COEF(1)+COEF(2)*P(1)+...,
4          23H+COEF(L)*P(1)*...P(L-1))
8          FORMAT(/10X,6HY(L,N,,I1,1H))
9          FORMAT(/16X,23HDET IN MIXED-RADIX FORM/10X,
1          35HFROM LEFT TO RIGHT, THE NUMBERS ARE,
2          24H COEF(1),...,COEF(L),AND/10X,
3          30HDET = COEF(1)+COEF(2)*P(1)+...,
4          23H+COEF(L)*P(1)*...P(L-1))
10         FORMAT(16X,20(I6))
11         FORMAT(/16X,21HY IN FIXED-RADIX FORM/10X,
1          35HFROM LEFT TO RIGHT, THE NUMBERS ARE,
2          24H COEF(1),...,COEF(L),AND/10X,
3          37HY(I,J) = COEF(1)*BASE**(L-1)+COEF(2)*,
4          23HBASE**(L-2)+...+COEF(L))
12        FORMAT(/16X,23HDET IN FIXED-RADIX FORM/10X,
1          35HFROM LEFT TO RIGHT, THE NUMBERS ARE,
2          24H COEF(1),...,COEF(L),AND/10X,
3          34HDET = COEF(1)*BASE**(L-1)+COEF(2)*,
4          21HBASE(N-2)+...+COEF(L)/16X,10(I6))
          RETURN
          END
DRIV0180
DRIV0190
DRIV0200
DRIV0210
DRIV0220
DRIV0230
DRIV0240
DRIV0250
DRIV0260
DRIV0270
DRIV0280
DRIV0290
DRIV0300
DRIV0310
DRIV0320
DRIV0330
DRIV0340
CONTINUEDRIV0350
DRIV0360
DRIV0370
DRIV0380
DRIV0390
DRIV0400
DRIV0410
DRIV0420
DRIV0430
DRIV0440
DRIV0450
DRIV0470
DRIV0480
DRIV0490
DRIV0500
DRIV0510
DRIV0520
DRIV0530
DRIV0540
DRIV0550
DRIV0560
DRIV0570
DRIV0580
DRIV0590
DRIV0600
CONTINUEDRIV0610
DRIV0620
DRIV0630
DRIV0640
DRIV0650
DRIV0660
DRIV0670
DRIV0680
DRIV0690
DRIV0700
DRIV0710
DRIV0720
DRIV0730
DRIV0740
DRIV0750
DRIV0760
DRIV0770
DRIV0780
DRIV0790
DRIV0800
DRIV0810
DRIV0820
DRIV0830
DRIV0840
DRIV0850
DRIV0860
DRIV0870
DRIV0880
DRIV0890
DRIV0900
DRIV0910

```

```

SUBROUTINE ESOLVE(A,B,AMOD,T,N,M,NM,DET,Y,MAXPRM,NOCOE,
1 RTEST,IER)
C
C*****
C THIS SUBROUTINE SOLVES EXACTLY THE SYSTEM OF LINEAR EQUATIONS*
C AX=B, WHERE A( ,N,N) AND B( ,N,M) ARE MATRICES WITH MULTIPLE-
C PRECISION INTEGER COEFFICIENTS EXPRESSED IN MIXED-RADIX FORM.*
C THE SUBROUTINE NORMALLY RETURNS INTEGER VALUES FOR
C DET=DETERMINANT(A) AND Y=A(ADJOINT)*B, ALSO IN MIXED-RADIX
C FORM. IF THE SOLUTION X IS REQUIRED, THE USER NEED ONLY
C COMPUTE X=Y/DET. (FOR CONVERSION OF INTEGERS FROM MIXED-
C RADIX TO FIXED-RADIX FORMS AND CONVERSELY, SUBROUTINES FRADIX*
C AND MRADIX CAN BE USED.)
C
C PRIME IS A LINEAR ARRAY CONTAINING 100 DISTINCT PRIME
C (INPUT) INTEGERS IN ASCENDING ORDER. THE PRIMES ARE CHOSEN
C AS LARGE AS POSSIBLE SUBJECT TO THE CONDITION THAT
C FOR ALL I AND J PRIME(I)*PRIME(J) DOES NOT OVERFLOW
C AN INTEGER WORD. THESE PRIMES ARE USED AS RADII FOR
C THE REPRESENTATION OF INTEGERS IN MIXED-RADIX FORM.
C IPRIME IS A LINEAR ARRAY OF INTEGERS SUCH THAT FOR EACH K,
C (INPUT) IPRIME(K)* PRIME(1)*PRIME(2)*...*PRIME(K-1) = 1,
C MODULO PRIME(K).
C
C A IS AN INTEGER MATRIX OF DIMENSION T BY N BY N. THE
C (INPUT) FIRST DIMENSION CONTAINS THE COEFFICIENTS OF THE
C MIXED-RADIX REPRESENTATION OF THE MULTIPLE-PRECISION
C COMPONENTS OF THE N BY N MATRIX A( ,N,N). THAT IS,
C
C A( ,I,J) = A(1,I,J)
C + A(2,I,J)*PRIME(1)
C
C .
C
C .
C
C + A(T,I,J)*PRIME(1)*...*PRIME(T-1),
C
C FOR I,J=1,2,...,N.
C
C B IS AN INTEGER MATRIX OF DIMENSION T BY N BY M WITH
C (INPUT) A SIMILAR NOTATIONAL CORRESPONDENCE MADE FOR A ABOVE.*
C AMOD IS AN N BY N+M MATRIX USED FOR TEMPORARY STORAGE OF
C (TEMP) THE AUGMENTED MATRIX (A,B) MODULO THE VARIOUS PRIMES
C PRIME(K). THIS MATRIX IS INCLUDED IN THE ARGUMENT
C LIST ONLY TO PERMIT ITS DIMENSIONS TO BE VARIABLE.
C
C T IS THE NUMBER OF RADII, PRIME(1),...,PRIME(T), USED
C (INPUT) TO REPRESENT EACH COMPONENT OF A( ,N,N) AND B( ,N,M)
C IN MIXED-RADIX FORM.
C
C N IS THE NUMBER OF EQUATIONS AND THE NUMBER OF UNKNOWN*
C (INPUT) IN THE SYSTEM. (I.E., N IS THE SIZE OF THE SECOND
C AND THIRD DIMENSIONS OF A.)
C
C M IS THE NUMBER OF RIGHT-HAND VECTORS FOR WHICH THE
C (INPUT) SYSTEM IS TO BE SOLVED. (I.E., M IS THE SIZE OF THE
C THIRD DIMENSION OF B.)
C
C NM IS EQUAL TO N+M.
C (INPUT)
C
C DET IS A VECTOR OF DIMENSION MAXPRM WHICH CONTAINS THE
C (OUTPUT) COEFFICIENTS OF THE MIXED-RADIX REPRESENTATION OF
C DETERMINANT(A)
C
C = DET(1)
C + DET(2)*PRIME(1)
C
C .
C
C .
C
C + DET(MAXPRM)*PRIME(1)*...*PRIME(MAXPRM-1).
C
C Y IS A MATRIX OF DIMENSION MAXPRM BY N BY M. THE FIRST*
C (OUTPUT) DIMENSION CONTAINS THE COEFFICIENTS OF THE MIXED-
C RADIX REPRESENTATION OF Y( ,N,M) = A(ADJOINT)*B:
C
C Y( ,I,J)
C = Y(1,I,J)
C + Y(2,I,J)*PRIME(1)
C
C .
C
C .
C
C + Y(MAXPRM,I,J)*PRIME(1)*...*PRIME(MAXPRM-1).
C
C MAXPRM IS THE MAXIMUM NUMBER OF RADII, PRIME(1),...,
C (INPUT) PRIME(MAXPRM), REQUIRED TO REPRESENT DETERMINANT(A)
C AND A(ADJOINT)*B IN MIXED-RADIX FORM. MAXPRM SHOULD
C BE CHOSEN SO THAT
C
C ABS(DETERMINANT(A)), ABS(Y( ,I,J))

```

```

ESOL0010
ESOL0020
ESOL0030
ESOL0040
ESOL0050
ESOL0060
ESOL0070
ESOL0080
ESOL0090
ESOL0100
ESOL0110
ESOL0120
ESOL0130
ESOL0140
ESOL0150
ESOL0160
ESOL0170
ESOL0180
ESOL0190
ESOL0200
ESOL0210
ESOL0220
ESOL0230
ESOL0240
ESOL0250
ESOL0260
ESOL0270
ESOL0280
ESOL0290
ESOL0300
ESOL0310
ESOL0320
ESOL0330
ESOL0340
ESOL0350
ESOL0360
ESOL0370
ESOL0380
ESOL0390
ESOL0400
ESOL0410
ESOL0420
ESOL0430
ESOL0440
ESOL0450
ESOL0460
ESOL0470
ESOL0480
ESOL0490
ESOL0500
ESOL0510
ESOL0520
ESOL0530
ESOL0540
ESOL0550
ESOL0560
ESOL0570
ESOL0580
ESOL0590
ESOL0600
ESOL0610
ESOL0620
ESOL0630
ESOL0640
ESOL0650
ESOL0660
ESOL0670
ESOL0680
ESOL0690
ESOL0700
ESOL0710
ESOL0720
ESOL0730
ESOL0740
ESOL0750
ESOL0760

```

```

C          < PRIME(1)*...*PRIME(MAXPRM)/2, I=1,2,...,N, *      ESOL0770
C          J=1,2,...,M. *      ESOL0780
C      CHOOSING MAXPRM = N*T + N*LOG(N)/(2*LOG(PRIME(1))) *      ESOL0790
C      WILL SUFFICE. (A TIGHTER BOUND CAN BE OBTAINED USING *      ESOL0800
C      SUBROUTINE SUBBND.) MAXPRM MUST BE LESS THAN 101. *      ESOL0810
C      NOCOEF IS THE NUMBER OF RADII ACTUALLY USED TO REPRESENT *      ESOL0820
C      C(OUTPUT) DET AND Y( ,I,J) IN MIXED-RADIX FORM. *      ESOL0830
C      NOCOEF <= MAXPRM. *      ESOL0840
C      RTEST IF RTEST = .FALSE. , MAXPRM COEFFICIENTS IN THE *      ESOL0850
C      (INPUT) MIXED-RADIX REPRESENTATION OF DET AND Y( ,N,M) *      ESOL0860
C      ARE COMPUTED. *      ESOL0870
C      IF RTEST = .TRUE. , THE ALGORITHM CONTINUES TO *      ESOL0880
C      ITERATE UNTIL TR (USUALLY TR=T, BUT SOMETIMES *      ESOL0890
C      TR=T+1) SUCCESSIVE ZERO COEFFICIENTS IN THE *      ESOL0900
C      MIXED-RADIX REPRESENTATION OF DET AND Y( ,N,M) *      ESOL0910
C      ARE ENCOUNTERED, THAT IS, WHENEVER FOR *      ESOL0920
C      K = NOCOEF+1,...,NOCOEF+TR (NOCOEF+TR <= MAXPRM), *      ESOL0930
C      DET(K) = 0 *      ESOL0940
C      Y(K,I,J) = 0, I=1,2,...,N, J=1,2,...,M. *      ESOL0950
C      IF NOCOEF > 0, THE PROGRAM TERMINATES WITH THE *      ESOL0960
C      KNOWLEDGE THAT *      ESOL0970
C      (1) A IS SINGULAR IF DET=0, OR *      ESOL0980
C      (2) Y( ,N,M)/DET IS THE SOLUTION TO AX=B IF *      ESOL0990
C      DET ≠ 0. (*WARNING* IN THIS CASE, A COMMON *      ESOL1000
C      FACTOR IN Y( ,N,M) AND DET MAY HAVE *      ESOL1010
C      BEEN IGNORED. THAT IS, DET MIGHT NO *      ESOL1020
C      LONGER BE THE DETERMINANT OF A.) *      ESOL1030
C      IF NOCOEF = 0, A IS PROBABLY SINGULAR, BUT THE *      ESOL1040
C      PROGRAM CONTINUES TO ITERATE UNTIL MAXPRM *      ESOL1050
C      COEFFICIENTS OF DET AND Y( ,N,M) HAVE BEEN *      ESOL1060
C      COMPUTED. *      ESOL1070
C      IER IS AN ERROR CODE WHICH IS *      ESOL1080
C      (OUTPUT) 0 IF Y( ,N,M)/DET IS THE SOLUTION TO AX=B. *      ESOL1090
C      (*WARNING* DET MIGHT NOT BE THE DETERMINANT OF A) *      ESOL1100
C      1 IF MAXPRM COEFFICIENTS IN THE MIXED-RADIX *      ESOL1110
C      REPRESENTATION OF DET AND Y( ,N,M) HAVE BEEN *      ESOL1120
C      COMPUTED. SINCE THE USER CHOOSES THE VALUE OF *      ESOL1130
C      MAXPRM, THE PROGRAM CANNOT GUARANTEE THAT *      ESOL1140
C      Y( ,N,M)/DET IS THE SOLUTION TO AX=B. *      ESOL1150
C      2 IF A IS SINGULAR. *      ESOL1160
C      3 IF A IS SINGULAR, OR DET AND Y( ,N,M) ARE NON-ZERO *      ESOL1170
C      MULTIPLES OF PRIME(1)*...*PRIME(MAXPRM). *      ESOL1180
C      4 IF THE INPUT PARAMETERS ARE INCORRECT. *      ESOL1190
C      *      ESOL1200
C***** *      ESOL1210
C      INTEGER T,TR,A(T,N,N),B(T,N,M),DET(MAXPRM),Y(MAXPRM,N,M), *      ESOL1220
C      1 AMOD(N,NM) *      ESOL1230
C      LOGICAL RTEST *      ESOL1240
C      INTEGER PT,TK,TT,TT1,DMOD,V1,V3,U1,U3,T1,T3,Q *      ESOL1250
C      INTEGER PRIME(100),IPRIME(100),P,P1,P2 *      ESOL1260
C      LOGICAL TEST *      ESOL1270
C      COMMON /PRIMEB/PRIME,IPRIME *      ESOL1280
C      IER=4 *      ESOL1290
C      IF (MAXPRM.GT.100 .OR. N.LT.2 .OR. M.LT.1) RETURN *      ESOL1300
C      IF (N+M.NE.NM) RETURN *      ESOL1310
C      N1 = N+1 *      ESOL1320
C      *      ESOL1330
C      *      ESOL1340
C      PROCEDURE ZERO-CRITERION ***** *      ESOL1350
C      DETERMINED IN THIS PROCEDURE IS THE NUMBER TR OF *      ESOL1360
C      CONSECUTIVE ZERO COEFFICIENTS IN THE MIXED-RADIX *      ESOL1370
C      REPRESENTATION OF DET AND Y REQUIRED TO GUARANTEE THAT *      ESOL1380
C      X=Y/DET IS THE SOLUTION OF AX=B. TR=T IF NORM INFINITY *      ESOL1390
C      OF THE MATRIX A AUGMENTED BY ANY COLUMN VECTOR OF B IS *      ESOL1400
C      BOUNDED BY 2*PRIME(1)*...*PRIME(T); OTHERWISE, TR=T+1. *      ESOL1410
C      *      ESOL1420
C      TR = T *      ESOL1430
C      IF (RTEST) *      ESOL1440
C      *      GOTO 1
C      *      GOTO 6
C      THEN *      ESOL1450
C      1 PT = 2*PRIME(T) *      ESOL1460
C      I = 0 *      ESOL1470
C      REPEAT ... FOR EACH ROW I *      ESOL1480
C      2 I = I+1 *      ESOL1490
C      IF (I.LE.N .AND. TR.EQ.T) *      ESOL1500
C      *      GOTO 3
C      *      GOTO 6

```



```

C          THEN
3          NORM = 0
            DO 4 J=1,N
4          NORM = NORM + IABS(A(T,I,J))
C          OD
            MAXB = 0
            DO 5 J=1,M
5          IF (MAXB.LT.IABS(B(T,I,J)))
1          MAXB = IABS(B(T,I,J))
C          OD
            NORM = NORM + MAXB
            IF (T.GT.1) NORM = NORM + N1
C          **HERE, USED IS THE FACT THAT FOR  $T \geq 1$ 
C          ** $(ABS(K(T)+1)*PRIME(1)*...*PRIME(T-1))$  IS
C          **A BOUND FOR ANY INTEGER GIVEN BY
C          ** $K(1)+K(2)*PRIME(1)+...+K(T)*PRIME(1)*$ 
C          ** $...*PRIME(T-1)$ .
            IF (NORM.GT.PT) TR = T+1
C          CONTINUE
                                                    GOTO 2
C          *
C          END ZERO--CRITERION *****
C          *
6          NOZERO = 0
          DO 64 ITER=1,MAXPRM
C          **THE SYSTEM AX=B IS SOLVED MODULO PRIME(ITER) FOR ALL
C          **ITER = 1,2,...,MAXPRM, OR IF RTEST IS TRUE UNTIL
C          **THE NUMBER OF CONSECUTIVE ZERO COEFFICIENTS (NOZERO)
C          **IN THE MIXED-RADIX REPRESENTATION OF DET AND Y IS
C          **EQUAL TO TR.
            P = PRIME(ITER)
            P1 = P-1
            P2 = P1/2 + 1
            IP = IPRIME(ITER)
C          *
C          PROCEDURE MAP *****
C          THIS PROCEDURE COMPUTES AMOD, THE AUGMENTED *
C          N BY N+M MATRIX (A,B) MODULO P. THAT IS, FOR *
C          I=1,2,...,N, HORNER'S RULE AND MODULO P ARITHMETIC *
C          IS USED TO COMPUTE *
C          AMOD(I,J) = A(1,I,J) + A(2,I,J)*PRIME(1) + *
C          ... + A(T,I,J)*PRIME(1)*...*PRIME(T-1), *
C          J=1,2,...,N,*
C          AMOD(I,J+N) = B(1,I,J) + B(2,I,J)*PRIME(1) + *
C          ... + B(T,I,J)*PRIME(1)*...*PRIME(T-1), *
C          J=1,2,...,M.*
C          *
            TT = MIN0(T,ITER)
            TT1 = TT-1
            IF (TT.EQ.1)
C          *
C          THEN
7          DO 10 I=1,N
            DO 8 J=1,N
8          AMOD(I,J) = A(1,I,J)
C          OD
            DO 9 J=1,M
            NJ = N+J
9          AMOD(I,NJ) = B(1,I,J)
C          OD
C          OD
10         CONTINUE
C          *
C          ELSE
11         DO 16 I=1,N
            DO 13 J=1,N
            NTEMP = A(TT,I,J)
            DO 12 K=1,TT1
            TK = TT - K
12         NTEMP = MOD(PRIME(TK)*NTEMP+A(TK,I,J),P)
C          OD
13         AMOD(I,J) = NTEMP
C          OD
            DO 15 J=1,M
            NTEMP = B(TT,I,J)
            DO 14 K=1,TT1
            NTEMP = MOD(PRIME(TK)*NTEMP+B(TK,I,J),P)
            DO 14 K=1,TT1

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      TK = TT - K
14      NTEMP = MOD(PRIME(TK)*NTEMP+B(TK,I,J),P)
C      OD
      NJ = N+J
15      AMOD(I,NJ) = NTEMP
C      OD
C      OD
16      CONTINUE
C      *
C      END MAP *****
C      *
C      PROCEDURE MSOLVE *****
C      THIS PROCEDURE SOLVES THE SYSTEM OF EQUATIONS *
C      AX=B MODULO P BY COMPUTING *
C      DMOD = DETERMINANT(A) MOD P *
C      AND *
C      A(ADJOINT)*B MOD P, *
C      WHICH IS TEMPORARILY STORED IN Y(ITER,I,J), *
C      I=1,2,...,N, J=1,2,...,M. *
C      *
C      THE GAUSSIAN ELIMINATION METHOD WITH PARTIAL *
C      PIVOTING IS USED TO REDUCE A TO UPPER ECHLON *
C      FORM. *
C      *
C      NSING ASSUMES THE VALUE *
C      0, IF A HAS RANK N, MOD P *
C      1, IF A HAS RANK N-1, MOD P *
C      2, IF A HAS RANK SMALLER THAN N-1, MOD P *
C      *
C      KRAM IS EQUAL TO N IF A IS NONSINGULAR MODULO P; *
C      OTHERWISE, KRAM IS THE INDEX OF THE FIRST *
C      COLUMN WHICH CONTAINS A ZERO PIVOT ELEMENT. *
C      *
17      DMOD = 1
      NSING = 0
      II = 0
C      REPEAT ... FORWARD ELIMINATION WITH II AS THE PIVOT
C      ... COLUMN
18      II = II+1
      IF (NSING.LT.2 .AND. II.LE.N)
C      THEN
19      I = II - NSING
C      **ELIMINATION PROCEEDS ON THE SUBMATRIX WITH
C      **ROWS I,I+1,...,N AND COLUMNS II,II+1,...,N+M
      II = I+1
      III = II+1
      K = I-1
C      REPEAT ... SEARCH FOR NON-ZERO ELEMENT IN
C      ... II'TH COLUMN.
20      K = K+1
      IF (K.LE.N)
C      THEN
21      IF (AMOD(K,II).NE.0)
C      THEN
C      c-----EXIT
C      CONTINUE
22      IF (K.LE.N)
C      THEN ... AMOD(K,II) IS THE PIVOT ELEMENT
23      IF (K.NE.I)
C      THEN ... INTERCHANGE ROWS I AND K
24      DO 25 J=II,NM
          NTEMP = AMOD(I,J)
          AMOD(I,J) = AMOD(K,J)
25      AMOD(K,J) = NTEMP
C      OD
C      DMOD = -DMOD
C      *
C      PROCEDURE INVERT *****
C      EUCLID'S EXTENDED ALGORITHM IS USED TO *
C      FIND THE INVERSE, VI, MODULO P OF *
C      AMOD(I,II). *
26      VI = 1

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```

      U1 = 0
      V3 = IABS(AMOD(I,II))
      U3 = P
      REPEAT
C      27      IF (V3.NE.1)
      GOTO 28
      GOTO 29
      THEN
C      28      Q = U3/V3
      T1 = U1 - Q*V1
      T3 = U3 - Q*V3
      U1 = V1
      U3 = V3
      V1 = T1
      V3 = T3
      GOTO 27
      CONTINUE
C      29      IF (AMOD(I,II).LT.0) V1 = -V1
      *
C      END INVERT *****
C      DMOD = MOD(DMOD*AMOD(I,II),P)
      DO 30 J=II, NM
C      30      AMOD(I,J) = MOD(V1*AMOD(I,J),P)
      OD
C      IF (I.LT.N)
      GOTO 31
      GOTO 34
      THEN
C      31      DO 32 K=II,N
      NTEMP = AMOD(K,II)
      DO 32 J=II, NM
C      32      AMOD(K,J) = MOD(AMOD(K,J) - NTEMP*
      1      AMOD(I,J),P)
      OD
C      OD
      GOTO 18
      ELSE ... COLUMN II IS A ZERO PIVOT COLUMN.
C      33      NSING = NSING + 1
      IF (NSING.LT.2) KRAM = II
      GOTO 18
      CONTINUE
C      34      IF (NSING.LT.2)
      GOTO 35
      GOTO 42
      THEN ... RANK OF A MODULO P IS N OR N-1; THEREFORE
      ... BACK SUBSTITUTE.
C      35      IF (NSING.EQ.0) KRAM = N
      KRAM1 = KRAM - 1
      KRAM2 = KRAM + 1
      DO 41 J=N1, NM
      NJ = J - N
      IF (KRAM.NE.N)
      GOTO 36
      GOTO 38
      THEN
C      36      DO 37 I=KRAM2,N
      Y(ITER,I,NJ) = 0
      OD
C      38      NTEMP = MOD(DMOD*AMOD(N,J),P)
      IF (NSING.EQ.1) NTEMP = NTEMP*(-1)**(N-KRAM)
      Y(ITER,KRAM,NJ) = NTEMP
      DO 40 II=1, KRAM1
      I = KRAM - II
      II = I+1
      NTEMP = 0
      DO 39 K=II, KRAM
C      39      NTEMP = MOD(NTEMP+AMOD(I,K)*
      1      Y(ITER,K,NJ), P)
      OD
      NTEMP = -NTEMP
      IF (NSING.EQ.0) NTEMP = MOD(NTEMP + DMOD*
      1      AMOD(I,J),P)
      Y(ITER,I,NJ) = NTEMP
      OD
      OD
      CONTINUE
      GOTO 44
      ELSE ... RANK OF A MODULO P IS LESS THAN N-1.
C      42      DO 43 J=1,M
      ESOL3060
      ESOL3070
      ESOL3080
      ESOL3090
      ESOL3100
      ESOL3110
      ESOL3120
      ESOL3130
      ESOL3140
      ESOL3150
      ESOL3160
      ESOL3170
      ESOL3180
      ESOL3190
      ESOL3200
      ESOL3210
      ESOL3220
      ESOL3230
      ESOL3240
      ESOL3250
      ESOL3260
      ESOL3270
      ESOL3280
      ESOL3290
      ESOL3300
      ESOL3310
      ESOL3320
      ESOL3330
      ESOL3340
      ESOL3350
      ESOL3360
      ESOL3370
      ESOL3380
      ESOL3390
      ESOL3400
      ESOL3410
      ESOL3420
      ESOL3430
      ESOL3440
      ESOL3450
      ESOL3460
      ESOL3470
      ESOL3480
      ESOL3490
      ESOL3500
      ESOL3510
      ESOL3520
      ESOL3530
      ESOL3540
      ESOL3550
      ESOL3560
      ESOL3570
      ESOL3580
      ESOL3590
      ESOL3600
      ESOL3610
      ESOL3620
      ESOL3630
      ESOL3640
      ESOL3650
      ESOL3660
      ESOL3670
      ESOL3680
      ESOL3690
      ESOL3700
      ESOL3710
      ESOL3720
      ESOL3730
      ESOL3740
      ESOL3750
      ESOL3760
      ESOL3770
      ESOL3780
      ESOL3790
      ESOL3800
      ESOL3810

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          DO 43 I=1,N
43          Y(ITER,I,J) = 0
C          OD
C          OD
44          IF (NSING.NE.0) DMOD = 0
C          *
C          END MSOLVE *****
C          PROCEDURE MIXED-RADIX *****
C          THIS PROCEDURE COMPUTES THE ITER'TH TERMS OF THE *
C          MIXED-RADIX REPRESENTATION OF Y( ,N,M) AND DET *
C          USING THE CHINESE REMAINDER THEOREM. *
C          *
C          TEST IS TRUE ONLY IF THE ITER'TH TERMS OF THE *
C          MIXED-RADIX REPRESENTATIONS OF DET AND Y ARE *
C          ALL ZERO. *
C          *
C          NOTE: IN FORTRAN, THE EVALUATION B=MOD(A,P) YIELDS *
C          AN INTEGER B SUCH THAT -P<B<P. IN THIS PROCEDURE, *
C          ANY SUCH B IS CONVERTED TO A SYMMETRIC RESIDUE C, *
C          -P<2C<P VIA THE EVALUATION C=B-(B/P2)*P. *
C          *
C          TEST = .TRUE.
C          ITER1 = ITER - 1
C          ITER2 = ITER - 2
C          DO 52 I=1,N
C          DO 52 J=1,M
C          IF (ITER.EQ.1)
C          THEN
45          MULT = Y(ITER,I,J)
C          GOTO 45 ESOL4090
C          GOTO 46 ESOL4100
C          ELSE
46          MULT = Y(ITER1,I,J)
C          IF (ITER.NE.2)
C          GOTO 47 ESOL4160
C          GOTO 49 ESOL4170
C          THEN
47          DO 48 LL=1,ITER2
C          L = ITER1 - LL
48          MULT = MOD(MULT*PRIME(L)+Y(L,I,J),P)
C          ESOL4210
C          OD
49          MULT = MOD(IP * MOD(Y(ITER,I,J)-MULT,P),P)
50          Y(ITER,I,J) = MULT - (MULT/P2)*P
50          IF (TEST)
C          GOTO 51 ESOL4250
C          GOTO 52 ESOL4260
C          THEN
51          IF (Y(ITER,I,J).NE.0) TEST = .FALSE.
52          ESOL4280
C          CONTINUEESOL4290
C          OD
C          OD
C          IF (ITER.EQ.1)
C          GOTO 53 ESOL4320
C          GOTO 54 ESOL4330
C          THEN
53          MULT = DMOD
C          ESOL4340
C          GOTO 58 ESOL4360
C          ELSE
54          MULT = DET(ITER1)
C          IF (ITER.NE.2)
C          GOTO 55 ESOL4390
C          GOTO 57 ESOL4400
C          THEN
55          DO 56 LL=1,ITER2
C          L = ITER1 - LL
56          MULT = MOD(MULT*PRIME(L) + DET(L),P)
C          ESOL4440
C          OD
57          MULT = MOD(IP * MOD(DMOD-MULT,P), P)
58          DET(ITER) = MULT - (MULT/P2)*P
58          IF (DET(ITER).NE.0) TEST = .FALSE.
C          ESOL4480
C          *
C          END MIXED-RADIX *****
C          *
C          IF (TEST)
C          GOTO 59 ESOL4520
C          GOTO 60 ESOL4530
C          THEN
59          NOZERO = NOZERO + 1
C          ESOL4540
C          ESOL4550
C          GOTO 61 ESOL4560
C          ELSE
C          ESOL4570

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60          NOZERO = 0
61          IF (RTEST .AND. ITER.GT.NOZERO .AND. NOZERO.GE.TR) GOTO 62
C          THEN ... NORMAL EXIT; FIRST, HOWEVER CHECK FOR ZERO
C          ... DETERMINANT.
62          IER = 0
          NOCOEF = ITER - NOZERO
          DO 63 K=1,NOCOE
            IF (DET(K).NE.0) RETURN
C          OD
63          IER = 2
          RETURN
64 CONTINUE
C
C          **MAXPRM COEFFICIENTS HAVE BEEN USED, ADNORMAL EXIT.
C
          NOCOEF = MAXPRM - NOZERO
          IER = 1
          IF (NOCOE.EQ.0) IER = 3
          RETURN
          END
          ESOL4580
          ESOL4590
          ESOL4600
          ESOL4610
          ESOL4620
          ESOL4630
          ESOL4640
          ESOL4650
          ESOL4660
          ESOL4670
          CONTINUEESOL4680
          ESOL4690
          ESOL4700
          ESOL4710
          ESOL4720
          ESOL4730
          ESOL4740
          ESOL4750
          ESOL4760
          ESOL4770
          ESOL4780
          ESOL4790

          BLOCK DATA
C          DATA0010
C          DATA0020
C*****
C          DATA0040
C          *****WARNING*****
C          DATA0050
C          DATA0060
C          DATA0070
C          DATA0080
C          DATA0090
C          DATA0100
C          DATA0110
C          DATA0120
C*****
C          DATA0130
C          DATA0140
C          DATA0150
          COMMON /PRIMEB/KPRIME(100),IPRIME(100)
          DATA KPRIME/
          45233,45247,45259,45263,45281,45289
          1,45293,45307,45317,45319,45329,45337,45341,45343,45361
          2,45377,45389,45403,45413,45427,45433,45439,45481,45491
          3,45497,45503,45523,45533,45541,45553,45557,45569,45587
          4,45589,45599,45613,45631,45641,45659,45667,45673,45677
          5,45691,45697,45707,45737,45751,45757,45763,45767,45779
          6,45817,45821,45823,45827,45833,45841,45853,45863,45869
          7,45887,45893,45943,45949,45953,45959,45971,45979,45989
          8,46021,46027,46049,46051,46061,46073,46091,46093,46099
          9,46103,46133,46141,46147,46153,46171,46181,46183,46187
          X,46199,46219,46229,46237,46261,46271,46273,46279,46301
          Y,46307,46309,46327,46337/
          DATA IPRIME/
          00000,42015,28577,01108,29342,16641
          1,10405,19447,26685,39525,14116,12753,32178,01043,08857
          2,27911,15049,07079,33425,00804,23175,23886,44779,41942
          3,10171,16606,10638,17371,27195,35827,42639,01829,24658
          4,09023,37958,30638,06339,41270,40538,10157,11783,00457
          5,32947,42170,17910,33474,20017,25086,36508,37444,35543
          6,06993,10326,16328,26765,42083,37223,30711,09408,06635
          7,38421,11397,32683,17333,34245,15748,35735,23492,19302
          8,20076,45620,44978,09864,14832,16092,19457,24045,44950
          9,32872,24309,15726,43057,37766,14046,41826,19946,41363
          X,23967,39791,29237,18085,12952,36850,02213,30023,34871
          Y,42667,40410,32615,46136/
          END
          DATA0020
          DATA0030
          DATA0040
          DATA0050
          DATA0060
          DATA0070
          DATA0080
          DATA0090
          DATA0100
          DATA0110
          DATA0120
          DATA0130
          DATA0140
          DATA0150
          DATA0160
          DATA0170
          DATA0180
          DATA0190
          DATA0200
          DATA0210
          DATA0220
          DATA0230
          DATA0240
          DATA0250
          DATA0260
          DATA0270
          DATA0280
          DATA0290
          DATA0300
          DATA0310
          DATA0320
          DATA0330
          DATA0340
          DATA0350
          DATA0360
          DATA0370
          DATA0380
          DATA0390
          DATA0400

          SUBROUTINE SUBBND(A,B,GAMMA,DELTA,S,SUM,T,N,M,NM,
          1 TWOTI,BOUND,NO)
C          SBND0010
C          SBND0020
C          SBND0030
C*****
C          SBND0040
C          SBND0050
C          SBND0060
C          SBND0070
C          SBND0080
C          SBND0090
C USING HADAMARD'S INEQUALITY, THIS SUBROUTINE COMPUTES THE
C MAXIMUM NUMBER OF PRIMES REQUIRED TO REPRESENT DETERMINANT(A)*
C AND A(ADJOINT)*B FOR THE SYSTEM OF LINEAR EQUATIONS AX=B,
C WHERE A( ,N,N) AND B( ,N,M) ARE MATRICES WITH MULTIPLE-

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C PRECISION INTEGER COEFFICIENTS EXPRESSED IN MIXED-RADIX FORM.* SBND0100
C * SBND0110
C PRIME IS A LINEAR ARRAY CONTAINING 100 DISTINCT PRIME * SBND0120
C (INPUT) INTEGERS IN ASCENDING ORDER. THESE PRIMES ARE THE * SBND0130
C RADII USED IN THE REPRESENTATION OF A AND B AND OF * SBND0140
C DETERMINANT(A) AND A(ADJOINT)*B.THE PRIMES ARE CHOSEN* SBND0150
C AS LARGE AS POSSIBLE SUBJECT TO THE CONDITION THAT * SBND0160
C FOR ALL I AND J PRIME(I)*PRIME(J) DOES NOT OVERFLOW * SBND0170
C AN INTEGER WORD. * SBND0180
C A IS AN INTEGER MATRIX OF DIMENSION T BY N BY N. THE * SBND0190
C (INPUT) FIRST DIMENSION CONTAINS THE COEFFICIENTS OF THE * SBND0200
C MIXED-RADIX REPRESENTATION OF THE MULTIPLE-PRECISION * SBND0210
C COMPONENTS OF THE N BY N MATRIX A( ,N,N). THAT IS, * SBND0220
C A( ,I,J) = A(1,I,J) * SBND0230
C + A(2,I,J)*PRIME(1) * SBND0240
C . * SBND0250
C . * SBND0260
C . * SBND0270
C + A(T,I,J)*PRIME(1)*...*PRIME(T-1), * SBND0280
C FOR I,J=1,2,...,N. * SBND0290
C B IS AN INTEGER MATRIX OF DIMENSION T BY N BY M WITH * SBND0300
C (INPUT) A SIMILAR NOTATIONAL CORRESPONDENCE MADE FOR A ABOVE.* SBND0310
C DELTA IS AN INTEGER ARRAY OF DIMENSION TWOT1 USED TO STORE * SBND0320
C (TEMP) THE COEFFICIENTS OF THE MULTIPLE-PRECISION INTEGERS * SBND0330
C A( ,I,J)**2 AND B( ,I,J)**2 FOR ANY PARTICULAR * SBND0340
C A( ,I,J) OR B( ,I,J). * SBND0350
C GAMMA IS AN INTEGER ARRAY OF DIMENSION TWOT1 USED TO STORE * SBND0360
C (TEMP) THE COEFFICIENTS OF THE MULTIPLE-PRECISION PARTIAL * SBND0370
C PRODUCTS OBTAINED WHEN A( ,I,J)**2 OR B( ,I,J)**2 ARE* SBND0380
C BEING COMPUTED. * SBND0390
C SUM IS AN INTEGER ARRAY OF DIMENSION TWOT1 USED TO STORE * SBND0400
C (TEMP) THE SUM OF (A( ,I,J)**2, I=1,...,N) OR THE SUM OF * SBND0410
C (B( ,I,J)**2 ,I=1,...,N) FOR ANY PARTICULAR COLUMN J.* SBND0420
C S IS A REAL ARRAY OF DIMENSION NM. FOR 1<=J<=N, S(J) * SBND0430
C (TEMP) CONTAINS A BOUND FOR THE LOGARITHM OF SUM(A( ,I,J)**2* * SBND0440
C ,I=1,...,N) AND FOR N+1<=J<=NM OF SUM(B( ,I,J)**2, * SBND0450
C I=1,...,N). * SBND0460
C T IS THE NUMBER OF RADII, PRIME(1),...,PRIME(T), USED * SBND0470
C (INPUT) TO REPRESENT EACH COMPONENT OF A( ,N,N) AND B( ,N,M) * SBND0480
C IN MIXED-RADIX FORM. * SBND0490
C N IS THE NUMBER OF EQUATIONS AND THE NUMBER OF UNKNOWN* SBND0500
C (INPUT) IN THE SYSTEM. (I.E., N IS THE SIZE OF THE SECOND * SBND0510
C AND THIRD DIMENSIONS OF A.) * SBND0520
C M IS THE NUMBER OF RIGHT-HAND VECTORS FOR WHICH THE * SBND0530
C (INPUT) SYSTEM IS TO BE SOLVED. (I.E., M IS THE SIZE OF THE * SBND0540
C THIRD DIMENSION OF B.) * SBND0550
C NM IS EQUAL TO N+M. * SBND0560
C (INPUT) * SBND0570
C TWOT1 IS EQUAL TO 2*T+1. * SBND0580
C (INPUT) * SBND0590
C NO IS A BOUND FOR THE NUMBER OF PRIMES REQUIRED TO SOLVE* SBND0600
C (OUTPUT) THE SYSTEM OF LINEAR EQUATIONS AX=B. THAT IS, NO IS * SBND0610
C A BOUND FOR THE NUMBER OF PRIMES REQUIRED TO * SBND0620
C REPRESENT DETERMINANT(A) AND A(ADJOINT)*B IN MIXED- * SBND0630
C RADIX FORM. * SBND0640
C BOUND IS EQUAL TO THE LOGARITHM OF THE PRODUCT OF THE FIRST* SBND0650
C (OUTPUT) NO PRIMES. * SBND0660
C * SBND0670
C ***** SBND0680
C INTEGER T,TWOT,TWOT1,A(T,N,N),B(T,N,M),DELTA(TWOT1), SBND0690
1 GAMMA(TWOT1),PRIME(100),SUM(TWOT1),CSUM,Q,PD SBND0700
REAL S(NM) SBND0710
INTEGER TK,TK1 SBND0720
COMMON /PRIMEB/ PRIME, IPRIME(100) SBND 730
TWOT = 2*T SBND0740
C PROCEDURE ABOUND***** SBND0750
C FOR EACH J, 1<=J<=N, THE SUM OF THE SQUARES OF THE * SBND0760
C ELEMENTS IN THE J-TH COLUMN OF A IS COMPUTED. THE * SBND0770
C SUM, WHICH IS A MIXED-RADIX, MULTIPLE-PRECISION * SBND0780
C INTEGER OF LENGTH AT MOST 2*T+1, IS STORED IN THE * SBND0790
C VECTOR SUM. THE LOGARITHM OF A BOUND FOR THE SUM IS * SBND0800
C THEN COMPUTED AND STORED IN S(J). * SBND0810
C * SBND0820
C DO 16 J=1,N SBND0830

```

```

C          **COMPUTE LOG (BOUND FOR (SUM OF A( ,I,J)**2,
C          **I=1,...,N)), AND STORE THE RESULTS IN S(J).
C          DO 1 K=1,TWOT1
1          SUM(K) = 0
C          OD
C          DO 9 I=1,N
C          PROCEDURE SQUARE*****
C          A( ,I,J)**2 IS COMPUTED. THE COEFFICIENTS OF *
C          THE MIXED-RADIX, MULTIPLE-PRECISION RESULT *
C          ARE STORED IN DELTA. *
C          *
C          DO 2 II=1,TWOT1
C          DELTA(II) = 0
2          GAMMA(II) = 0
C          OD
C          Q = 0
C          K = 0
C          L = T-1
C          REPEAT
3          IF (K.LE.L)
C          THEN
C          **COMPUTE PARTIAL PRODUCT AND STORE IN
C          **GAMMA.
4          Q = 0
C          TK = T-K
C          DO 5 II=1,T
C          PD = A(TK,I,J)*A(II,I,J) + Q
C          Q = PD/PRIME(II)
5          GAMMA(II) = PD - Q*PRIME(II)
C          OD
C          GAMMA(T+1) = Q
C          **ADD PARTIAL PRODUCT TO DELTA.
C          Q = 0
C          L1 = T+K
C          TK1 = T+K+1
C          DO 6 II=1,L1
C          PD = PRIME(TK)*DELTA(II) + GAMMA(II)
1          + Q
C          Q = PD/PRIME(II)
6          DELTA(II) = PD - Q*PRIME(II)
C          OD
C          DELTA(TK1) = Q
C          K = K+1
C          GOTO 3
C          CONTINUE
C          *****
C          END SQUARE*****
C          **ACCUMULATE THE SUM OF A( ,I,J)**2.
7          Q = 0
C          DO 8 II=1,TWOT1
C          CSUM = DELTA(II)+SUM(II)+Q
C          Q = CSUM/PRIME(II)
8          SUM(II) = CSUM - Q*PRIME(II)
C          OD
C          OD
9          CONTINUE
C          **CALCULATE ABOUND FOR THE MULTIPLE-PRECISION
C          **INTEGER IN SUM AND STORE THE LOG OF THIS
C          **BOUND IN S(J).
C          K = TWOT1
C          REPEAT...FIND OUT THE LENGTH K OF THE MULTIPLE-
C          ...PRECISION NUMBER IN SUM.
10         IF (K.GT.1.AND.SUM(K).EQ.0)
C          THEN
11         K = K - 1
C          GOTO 10
C          CONTINUE
12         S(J) = 0.0
C          IF (K.GT.1)
C          GOTO 13
C          GOTO 15

```

```

C          THEN                                SBND1600
13          L = K-1                            SBND1610
          DO 14 II=1,L                        SBND1620
14          S(J) = S(J) +ALOG(FLOAT(PRIME(II))) SBND1630
C          OD                                  SBND1640
          S(J) = S(J)+ALOG(FLOAT(SUM(K)+1))    SBND1650
C                                               GOTO 16 SBND1660
C          ELSE                                SBND1670
15          IF (SUM(K).GT.0) S(J) = ALOG(FLOAT(SUM(K))) SBND1680
C          OD                                  SBND1690
16          CONTINUESBND1700
C          * SBND1710
C          END ABOUND*****                    SBND1720
C          SBND1730
C          SBND1740
C          PROCEDURE BBOUND*****              SBND1750
C          FOR EACH J, 1<=J<=M, THE SUM OF THE SQUARES OF THE * SBND1760
C          ELEMENTS IN THE J-TH COLUMN OF B IS COMPUTED. THE * SBND1770
C          SUM, WHICH IS A MIXED-RADIX, MULTIPLE PRECISION * SBND1780
C          INTEGER OF LENGTH AT MOST 2*T+1, IS STORED IN THE * SBND1790
C          VECTOR SUM. THE LOGARITHM OF A BOUND FOR THE SUM IS * SBND1800
C          THEN COMPUTED AND STORED IN S(N+J). * SBND1810
C          * SBND1820
C          NOTE: THE COMPUTATIONAL PROCEDURE IS THE SAME AS FOR * SBND1830
C          A ABOVE. * SBND1840
C          * SBND1850
          DO 32 J=1,M                          SBND1860
          DO 17 K=1,TWOT1                      SBND1870
17          SUM(K) = 0                         SBND1880
C          OD                                  SBND1890
          DO 25 I=1,N                          SBND1900
          DO 18 II=1,TWOT1                    SBND1910
18          DELTA(II) = 0                     SBND1920
          GAMMA(II) = 0                       SBND1930
C          OD                                  SBND1940
          Q = 0                                SBND1950
          K = 0                                SBND1960
          L = T-1                              SBND1970
C          REPEAT                              SBND1980
19          IF (K.LE.L)                       GOTO 20 SBND1990
C                                               GOTO 23 SBND2000
C          THEN                                SBND2010
20          Q = 0                              SBND2020
          TK = T-K                             SBND2030
          DO 21 II=1,T                        SBND2040
          PD = B(TK,I,J)*B(II,I,J) + Q       SBND2050
          Q = PD/PRIME(II)                   SBND2060
21          GAMMA(II) = PD - Q*PRIME(II)     SBND2070
C          OD                                  SBND2080
          GAMMA(T+1) = Q                      SBND2090
          Q = 0                                SBND2100
          L1 = T+K                            SBND2110
          TK1 = T+K+1                        SBND2120
          DO 22 II=1,L1                      SBND2130
          PD = PRIME(TK)*DELTA(II) + GAMMA(II)+Q SBND2140
          Q = PD/PRIME(II)                   SBND2150
22          DELTA(II) = PD - Q*PRIME(II)     SBND2160
C          OD                                  SBND2170
          DELTA(TK1) = Q                      SBND2180
          K = K+1                             SBND2190
C                                               GOTO 19 SBND2200
C          CONTINUE                            SBND2210
23          Q = 0                              SBND2220
          DO 24 II=1,TWOT1                    SBND2230
          CSUM = DELTA(II)+SUM(II)+Q         SBND2240
          Q = CSUM/PRIME(II)                 SBND2250
24          SUM(II) = CSUM - Q*PRIME(II)     SBND2260
C          OD                                  SBND2270
C          OD                                  SBND2280
25          CONTINUESBND2290
          K = TWOT1                           SBND2300
C          REPEAT                              SBND2310
26          IF (K.GT.1.AND.SUM(K).EQ.0)     GOTO 27 SBND2320
C                                               GOTO 28 SBND2330
C          THEN                                SBND2340
27          K = K - 1                         SBND2350

```



```

C          CONTINUE
28      NJ = N+J
        S(NJ) = 0.0
        IF (K.GT.1)
C          THEN
29          L = K-1
            DO 30 II=1,L
30              S(NJ) = S(NJ) +ALOG(FLOAT(PRIME(II)))
C          OD
            S(NJ) = S(NJ)+ALOG(FLOAT(SUM(K)+1))
C          ELSE
31          IF (SUM(K).GT.0) S(J) = ALOG(FLOAT(SUM(K)))
C          OD
32          CONTINUE
C          *
C          END BBOUND*****
C
C          **FIND MIN(S(I)), I=1,...,N.
            SMIN = S(1)
            DO 33 I=1,N
33              IF (SMIN.GT.S(I)) SMIN = S(I)
C          OD
C          **FIND MAX(S(I)), I=N+1,...,NM.
            SMAX = S(N+1)
            L = N+1
            DO 34 I=L,NM
34              IF (SMAX.LT.S(I)) SMAX = S(I)
C          OD
C          **IF SMIN 1/2= SMAX, THEN BOUND = (SUM OF S(I))/2+ LOG(2.0)
C          **ELSE BOUND = ((SUM OF S(I))+SMAX-SMIN)/2+LOG(2.0),
C          **WHERE THE SUM IS TAKEN OVER I=1,...,N.
            BOUND = 0.0
            DO 35 I=1,N
35              BOUND = BOUND+S(I)
C          OD
            IF (SMIN.GE.SMAX)
C          THEN
36              BOUND = BOUND/2.0 + ALOG(2.0)
C          ELSE
37              BOUND = (BOUND+SMAX-SMIN)/2.0 + ALOG(2.0)
C          **CALCULATE NUMBER OF PRIMES REQUIRED.
            NO = 0
            SUMLOG = 0.0
            REPEAT
39              NO = NO+1
                IF (SUMLOG.GE.BOUND)
C          THEN
C          c-----EXIT
                SUMLOG = SUMLOG + ALOG(FLOAT(PRIME(NO)))
C          CONTINUE
40      RETURN
        END

SUBROUTINE MRADIX(A,N,C,LMAX,L,B,IER)
C
C*****
C GIVEN THE FIXED-RADIX INTEGER
C
C      A(1)*B**(N-1) + ... + A(N-1)*B + A(N),
C
C WHERE ABS(A(I)) < B, I=1,...,N,
C THIS SUBROUTINE COMPUTES (FOR SOME L <= LMAX) THE COEFFICIENTS*
C C(1),C(2),...,C(L) IN ITS SYMMETRIC MIXED-RADIX
C REPRESENTATION:
C
C      C(1)+C(2)*P(1) + ... + C(L)*P(1)*P(2)*...*P(L-1).
C
C LMAX MUST BE GIVEN SUCH THAT
C

```

```

GOTO 26 SBND2360
SBND2370
SBND2380
SBND2390
GOTO 29 SBND2400
GOTO 31 SBND2410
SBND2420
SBND2430
SBND2440
SBND2450
SBND2460
SBND2470
GOTO 32 SBND2480
SBND2490
SBND2500
SBND2510
CONTINUESBND2520
* SBND2530
SBND2540
SBND2550
SBND2560
SBND2570
SBND2580
SBND2590
SBND2600
SBND2610
SBND2620
SBND2630
SBND2640
SBND2650
SBND2660
SBND2670
SBND2680
SBND2690
SBND2700
SBND2710
SBND2720
SBND2730
SBND2740
GOTO 36 SBND2750
GOTO 37 SBND2760
SBND2770
SBND2780
GOTO 38 SBND2790
SBND2800
SBND2810
SBND2820
SBND2830
SBND2840
SBND2850
SBND2860
GOTO 40 SBND2870
SBND2880
SBND2890
SBND2900
GOTO 39 SBND2910
SBND2920
SBND2930
SBND2940

```

```

MRAD0010
MRAD0020
MRAD0030
MRAD0040
MRAD0050
MRAD0060
MRAD0070
MRAD0080
MRAD0090
MRAD0100
MRAD0110
MRAD0120
MRAD0130
MRAD0140
MRAD0150
MRAD0160

```

```

C   LMAX >= (N * LOG B + LOG 2) / LOG(P(1)).          *   MRAD0170
C                                                     *   MRAD0180
C   IER IS AN ERROR CODE WHICH IS 1 IF THE DIMENSION PARAMETER N,*
C   LMAX ARE INCORRECT, AND 0 OTHERWISE.              *   MRAD0190
C                                                     *   MRAD0200
C                                                     *   MRAD0210
C               ***** WARNING *****              *   MRAD0220
C                                                     *   MRAD0230
C   THIS SUBROUTINE ASSUMES THAT FOR ALL I,J           *   MRAD0240
C                                                     *   MRAD0250
C   P(I)*P(J) AND P(I)*B                             *   MRAD0260
C                                                     *   MRAD0270
C   DO NOT OVERFLOW A COMPUTER WORD.                 *   MRAD0280
C                                                     *   MRAD0290
C*****                                              *   MRAD0300
C                                                     *   MRAD0310
C   INTEGER A(N),C(LMAX),P(100),Q,QTEMP,PP,B,P1      *   MRAD0320
COMMON /PRIMEB/ P, IPRIME(100)                      *   MRAD 330
IER = 1                                              *   MRAD0340
IF (N.LT.1 .OR. LMAX.LT.1) RETURN                  *   MRAD0350
L = 1                                              *   MRAD0360
C(1) = 0                                           *   MRAD0370
DO 4 I=1,N                                         *   MRAD0380
C   **AT THIS STAGE C(1)+C(2)*P(1)+...+C(L)*P(1)*P(2)*...*
C   **P(L-1) IS THE MIXED-RADIX REPRESENTATION OF
C   **A(1)*B**(I-2)+A(2)*B**(I-3)+...+A(I-1).
C   Q = A(I)                                       *   MRAD0390
DO 1 J=1,L                                         *   MRAD0400
C   **COMPUTE THE FIRST L COEFFICIENTS OF THE MIXED-
C   **RADIX REPRESENTATION OF A(I)+B*(C(1)+C(2)*P(1)+
C   **...+C(L)*P(1)*P(2)*...*P(L-1)).
C   PP = P(J)                                       *   MRAD0410
QTEMP = Q + B*C(J)                                *   MRAD0420
Q = QTEMP/PP                                       *   MRAD0430
C(J) = QTEMP - Q*PP                                *   MRAD0440
1   OD                                             *   MRAD0450
C   REPEAT ... CONVERT C(1)+C(2)*P(1)+...+C(L)*P(1)*
C   ... P(2)*...*P(L-1)+Q*P(1)*...*P(L)
C   ... TO MIXED-RADIX FORM.
C   2   IF (Q .NE.0)                                *   MRAD0460
C                                                     *   MRAD0470
C   THEN                                           *   MRAD0480
C   3   L = L+1                                     *   MRAD0490
C   IF (L.GT.LMAX) RETURN                          *   MRAD0500
C   PP = P(L)                                       *   MRAD0510
C   QTEMP = Q/PP                                    *   MRAD0520
C   C(L) = Q - QTEMP*PP                             *   MRAD0530
C   Q = QTEMP                                       *   MRAD0540
C                                                     *   MRAD0550
C   CONTINUE                                       *   MRAD0560
C   OD                                             *   MRAD0570
C   4   CONTINUEMRAD0670
C   **CONVERT TO SYMMETRIC MIXED-RADIX FORM.
C   Q = 0
C   DO 5 I=1,L
C   C(I) = C(I)+Q
C   P1 = (P(I)+1)/2
C   Q = C(I)/P1
C   5   C(I) = C(I) - Q*P(I)
C   OD
C   IF (Q.NE.0)
C   6   THEN
C   L = L+1
C   IF (L.GT.LMAX) RETURN
C   C(L) = Q
C   7   IER = 0
RETURN
END
*   MRAD0680
*   MRAD0690
*   MRAD0700
*   MRAD0710
*   MRAD0720
*   MRAD0730
*   MRAD0740
*   MRAD0750
*   MRAD0760
*   MRAD0770
*   MRAD0780
*   MRAD0790
*   MRAD0800
*   MRAD0810
*   MRAD0820
*   MRAD0830
*   MRAD0840
SUBROUTINE FRADIX(C,N,A,LMAX,L,B,IER)
C
C*****
C   GIVEN THE MIXED-RADIX INTEGER
C
C   C(1)+C(2)*P(1)+ ... +C(N)*P(1)*...*P(N-1),
*   FRAD0010
*   FRAD0020
*   FRAD0030
*   FRAD0040
*   FRAD0050
*   FRAD0060

```

```

C                                     *      FRAD0070
C WHERE ABS(C(I)) < (P(I)+1)/2, I=1,...,N, *      FRAD0080
C THIS SUBROUTINE COMPUTES(FOR SOME L <= LMAX) THE COEFFICIENTS*      FRAD0090
C A(1),A(2),...,A(L) IN ITS FIXED-RADIX REPRESENTATION: *      FRAD0100
C                                     *      FRAD0110
C   A(1)*B**(L-1)+A(2)*B*(L-2)+...+A(L), *      FRAD0120
C                                     *      FRAD0130
C WHERE ABS(A(I)) < B. *      FRAD0140
C                                     *      FRAD0150
C LMAX MUST BE GIVEN SUCH THAT *      FRAD0160
C                                     *      FRAD0170
C   LMAX >= (N*LOG(P(N)) - LOG 2) / LOG B. *      FRAD0180
C                                     *      FRAD0190
C IER IS AN ERROR CODE WHICH IS 1 IF THE DIMENSION PARAMETER N,*      FRAD0200
C LMAX ARE INCORRECT, AND 0 OTHERWISE. *      FRAD0210
C                                     *      FRAD0220
C           ***** WARNING ***** *      FRAD0230
C                                     *      FRAD0240
C THIS SUBROUTINE ASSUMES THAT FOR ALL I *      FRAD0250
C                                     *      FRAD0260
C   P(I)*B *      FRAD0270
C                                     *      FRAD0280
C DOES NOT OVERFLOW A COMPUTER WORD. *      FRAD0290
C                                     *      FRAD0300
C***** *      FRAD0310
C                                     *      FRAD0320
C   INTEGER C(N),A(LMAX),P(100),Q,QTEMP,PP,B *      FRAD0330
C   COMMON /PRIMEB/ P, IPRIME(100) *      FRAD 340
C   IER = 1 *      FRAD0350
C   IF (N.LT.1 .OR. LMAX.LT.1) RETURN *      FRAD0360
C   L = 0 *      FRAD0370
C   DO 5 I=1,N *      FRAD0380
C     **AT THIS STAGE A(1)+A(2)*B+...+A(L)*B**(L-1) IS THE *      FRAD0390
C     **FIXED-RADIX REPRESENTATION OF C(N-I+2)+C(N-I+3)* *      FRAD0400
C     **P(N-I+2)+...+C(N)*P(N-I+2)*...*P(N-1). *      FRAD0410
C     NI1 = N-I+1 *      FRAD0420
C     PP = P(NI1) *      FRAD0430
C     Q = C(NI1) *      FRAD0440
C     IF (L.GT.0) *      FRAD0450
C                                     GOTO 1
C                                     GOTO 3
C     THEN ... COMPUTE THE FIRST L COEFFICIENTS OF THE *      FRAD0470
C     ... FIXED-RADIX REPRESENTATION OF C(N-I+1)+ *      FRAD0480
C     ... P(N-I+1)*(A(1)+A(2)*B+...+A(L)*B**(L-1)). *      FRAD0490
C   1   DO 2 J=1,L *      FRAD0500
C       QTEMP = A(J)*PP + Q *      FRAD0510
C       Q = QTEMP/B *      FRAD0520
C   2   A(J) = QTEMP - Q*B *      FRAD0530
C       OD *      FRAD0540
C   REPEAT ... CONVERT A(1)+A(2)*B+...+A(L)*B**(L-1) *      FRAD0550
C   ... +Q*B**L TO FIXED-RADIX FORM. *      FRAD0560
C   3   IF (Q.NE.0) *      FRAD0570
C                                     GOTO 4
C                                     GOTO 5
C   THEN *      FRAD0590
C   4   L = L+1 *      FRAD0600
C       IF (L.GT.LMAX) RETURN *      FRAD0610
C       QTEMP = Q/B *      FRAD0620
C       A(L) = Q - QTEMP*B *      FRAD0630
C       Q = QTEMP *      FRAD0640
C                                     GOTO 3
C   CONTINUE *      FRAD0660
C   5   CONTINUEFRAD0670
C   OD *      FRAD0680
C   IF (L.GE.2) *      FRAD0690
C                                     GOTO 6
C                                     GOTO 8
C   THEN ... REORDER THE COEFFICIENTS. *      FRAD0710
C   6   L2 = L/2 *      FRAD0720
C       DO 7 I=1,L2 *      FRAD0730
C         LI1 = L-I+1 *      FRAD0740
C         QTEMP = A(I) *      FRAD0750
C         A(I) = A(LI1) *      FRAD0760
C         A(LI1) = QTEMP *      FRAD0770
C   7   OD *      FRAD0780
C   8   IER = 0 *      FRAD0790
C       RETURN *      FRAD0800
C       END *      FRAD0810

```


ALGORITHM 523

CONVEX, A New Convex Hull Algorithm for Planar Sets [Z]

WILLIAM F. EDDY
Carnegie-Mellon University

Key Words and Phrases: convex hull, QUICKERSORT, partitioning, sorting
CR Categories: 5.30, 5.31
Language: Fortran

DESCRIPTION

An algorithm, CONVEX, that determines which points of a planar set are vertices of the convex hull of the set is presented here. A detailed explanation of its operation and a report of a small sampling experiment are given in [1].

REFERENCES

1. EDDY, W.F., A new convex hull algorithm for planar sets. *ACM Trans. Math. Software* 3, 4 (Dec. 1977), 398-403.

ALGORITHM

```

COMMON NCOUNT
DIMENSION XX(2,25),IN(25),IH(25)
DIMENSION X(25),Y(25)
INTEGER IWORK(50)
INTEGER IL(50)
DATA IWORK/50*0/
NCOUNT=0
C NCOUNT IS TOTAL NUMBER OF POINTS PASSED TO SPLIT
READ(5,1)N
1  FORMAT(I5)
   WRITE(6,1)N
   N1=N+1
   DO 2 I=1,N
     J=N1-I
2  IN(J)=I
C ARRAY IN CONTAINS INDICES 1-N IN REVERSE ORDER
DO 3 I=1,N
3  READ(5,4)XX(1,I),XX(2,I)
4  FORMAT(2F10.5)
   DO 5 I=1,N
     J=IN(I)
5  WRITE(6,4)XX(1,J),XX(2,J)
   DO 10 M=4,N
     CALL CONVEX(N,XX,M,IN,IWORK,IWORK(N+1),IH,NHULL,IL)
   IK=IL(1)

```

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      DO 6 I=1,NHULL
        J=IH(IK)
        X(I)=XX(1,J)
        Y(I)=XX(2,J)
6       IK=IL(IK)
        WRITE(6,7)M,NHULL,NCOUNT
7       FORMAT(12H0SAMPLE SIZE ,I5,9H VERTICES ,I5,6H SPLIT ,I5)
        DO 8 I=1,NHULL
8       WRITE(6,9)X(I),Y(I)
9       FORMAT(1X,2F10.5)
10      CONTINUE
        STOP
        END

```

```

      SUBROUTINE SPLIT(N,X,M,IN,II,JJ,S,IABV,NA,MAXA,IBEL,
1  NB,MAXB)
C THIS SUBROUTINE TAKES THE M POINTS OF ARRAY X WHOSE
C SUBSCRIPTS ARE IN ARRAY IN AND PARTITIONS THEM BY THE
C LINE JOINING THE TWO POINTS IN ARRAY X WHOSE SUBSCRIPTS
C ARE II AND JJ. THE SUBSCRIPTS OF THE POINTS ABOVE THE
C LINE ARE PUT INTO ARRAY IABV, AND THE SUBSCRIPTS OF THE
C POINTS BELOW ARE PUT INTO ARRAY IBEL. NA AND NB ARE,
C RESPECTIVELY, THE NUMBER OF POINTS ABOVE THE LINE AND THE
C NUMBER BELOW. MAXA AND MAXB ARE THE SUBSCRIPTS FOR ARRAY
C X OF THE POINT FURTHEST ABOVE THE LINE AND THE POINT
C FURTHEST BELOW, RESPECTIVELY. IF EITHER SUBSET IS NULL
C THE CORRESPONDING SUBSCRIPT (MAXA OR MAXB) IS SET TO ZERO
C FORMAL PARAMETERS
C INPUT
C N   INTEGER          TOTAL NUMBER OF DATA POINTS
C X   REAL ARRAY (2,N) (X,Y) CO-ORDINATES OF THE DATA
C M   INTEGER          NUMBER OF POINTS IN INPUT SUBSET
C IN  INTEGER ARRAY (M) SUBSCRIPTS FOR ARRAY X OF THE
C                               POINTS IN THE INPUT SUBSET
C II  INTEGER          SUBSCRIPT FOR ARRAY X OF ONE POINT
C                               ON THE PARTITIONING LINE
C JJ  INTEGER          SUBSCRIPT FOR ARRAY X OF ANOTHER
C                               POINT ON THE PARTITIONING LINE
C S   INTEGER          SWITCH TO DETERMINE OUTPUT. REFER
C                               TO COMMENTS BELOW
C OUTPUT
C IABV INTEGER ARRAY (M) SUBSCRIPTS FOR ARRAY X OF THE
C                               POINTS ABOVE THE PARTITIONING LINE
C NA  INTEGER          NUMBER OF ELEMENTS IN IABV
C MAXA INTEGER         SUBSCRIPT FOR ARRAY X OF POINT
C                               FURTHEST ABOVE THE LINE. SET TO
C                               ZERO IF NA IS ZERO
C IBEL INTEGER ARRAY (M) SUBSCRIPTS FOR ARRAY X OF THE
C                               POINTS BELOW THE PARTITIONING LINE
C NB  INTEGER          NUMBER OF ELEMENTS IN IBEL
C MAXB INTEGER         SUBSCRIPT FOR ARRAY X OF PCINT
C                               FURTHEST BELOW THE LINE. SET TO
C                               ZERO IF NB IS ZERO
      DIMENSION X(2,N)
      DIMENSION IN(M),IABV(M),IBEL(M)
      INTEGER S
C IF S = 2 DONT SAVE IBEL,NB,MAXB.
C IF S =-2 DONT SAVE IABV,NA,MAXA.
C OTHERWISE SAVE EVERYTHING
C IF S IS POSITIVE THE ARRAY BEING PARTITIONED IS ABOVE
C THE INITIAL PARTITIONING LINE. IF IT IS NEGATIVE, THEN
C THE SET OF POINTS IS BELOW.
      LOGICAL T
      T=.FALSE.
C CHECK TO SEE IF THE LINE IS VERTICAL
      IF(X(1,JJ).NE.X(1,II))GOTO 1
      XT=X(1,II)
      DIR=SIGN(1.,X(2,JJ)-X(2,II))*SIGN(1.,FLOAT(S))
      T=.TRUE.
      GOTO 2
1     A=(X(2,JJ)-X(2,II))/(X(1,JJ)-X(1,II))
      B=X(2,II)-A*X(1,II)
2     UP=0.
      NA=0
      MAXA=0
      DOWN=0.

```

```

        NB=0
        MAXB=0
        DO 6 I=1,M
            IS=IN(I)
            IF(T)GOTO 3
            Z=X(2,IS)-A*X(1,IS)-B
            GOTO 4
3         Z=DIR*(X(1,IS)-XT)
4         IF(Z.LE.0.)GOTO 5
C THE POINT IS ABOVE THE LINE
            IF(S.EQ.-2)GOTO 6
            NA=NA+1
            IABV(NA)=IS
            IF(Z.LT.UP)GOTO 6
            UP=Z
            MAXA=NA
            GOTO 6
5         IF(S.EQ.2)GOTO 6
            IF(Z.GE.0.)GOTO 6
C THE POINT IS BELOW THE LINE
            NB=NB+1
            IBEL(NB)=IS
            IF(Z.GT.DOWN)GOTO 6
            DOWN=Z
            MAXB=NB
6         CONTINUE
        RETURN
        END

```

```

        SUBROUTINE CONVEX(N,X,M,IN,IA,IB,IH,NH,IL)
C THIS SUBROUTINE DETERMINES WHICH OF THE M POINTS OF ARRAY
C X WHOSE SUBSCRIPTS ARE IN ARRAY IN ARE VERTICES OF THE
C MINIMUM AREA CONVEX POLYGON CONTAINING THE M POINTS. THE
C SUBSCRIPTS OF THE VERTICES ARE PLACED IN ARRAY IH IN THE
C ORDER THEY ARE FOUND. NH IS THE NUMBER OF ELEMENTS IN
C ARRAY IH AND ARRAY IL. ARRAY IL IS A LINKED LIST GIVING
C THE ORDER OF THE ELEMENTS OF ARRAY IH IN A COUNTER
C CLOCKWISE DIRECTION. THIS ALGORITHM CORRESPONDS TO A
C PREORDER TRAVERSAL OF A CERTAIN BINARY TREE. EACH VERTEX
C OF THE BINARY TREE REPRESENTS A SUBSET OF THE M POINTS.
C AT EACH STEP THE SUBSET OF POINTS CORRESPONDING TO THE
C CURRENT VERTEX OF THE TREE IS PARTITIONED BY A LINE
C JOINING TWO VERTICES OF THE CONVEX POLYGON. THE LEFT SON
C VERTEX IN THE BINARY TREE REPRESENTS THE SUBSET OF POINTS
C ABOVE THE PARTITIONING LINE AND THE RIGHT SON VERTEX, THE
C SUBSET BELOW THE LINE. THE LEAVES OF THE TREE REPRESENT
C EITHER NULL SUBSETS OR SUBSETS INSIDE A TRIANGLE WHOSE
C VERTICES COINCIDE WITH VERTICES OF THE CONVEX POLYGON.
C FORMAL PARAMETERS
C INPUT
C N INTEGER TOTAL NUMBER OF DATA POINTS
C X REAL ARRAY (2,N) (X,Y) CO-ORDINATES OF THE DATA
C M INTEGER NUMBER OF POINTS IN THE INPUT SUBSET
C IN INTEGER ARRAY (M) SUBSCRIPTS FOR ARRAY X OF THE POINTS
C IN THE INPUT SUBSET
C WORK AREA
C IA INTEGER ARRAY (M) SUBSCRIPTS FOR ARRAY X OF LEFT SON
C SUBSETS. SEE COMMENTS AFTER DIMENSION
C STATEMENTS
C IB INTEGER ARRAY (M) SUBSCRIPTS FOR ARRAY X OF RIGHT SON
C SUBSETS
C OUTPUT
C IH INTEGER ARRAY (M) SUBSCRIPTS FOR ARRAY X OF THE
C VERTICES OF THE CONVEX HULL
C NH INTEGER NUMBER OF ELEMENTS IN ARRAY IH AND
C ARRAY IL. SAME AS NUMBER OF VERTICES
C OF THE CONVEX POLYGON
C IL INTEGER ARRAY (M) A LINKED LIST GIVING IN ORDER IN A
C COUNTER-CLOCKWISE DIRECTION THE
C ELEMENTS OF ARRAY IH
        DIMENSION X(2,N)
        DIMENSION IN(M),IA(M),IB(M),IH(M),IL(M)
C THE UPPER END OF ARRAY IA IS USED TO STORE TEMPORARILY
C THE SIZES OF THE SUBSETS WHICH CORRESPOND TO RIGHT SON
C VERTICES, WHILE TRAVERSING DOWN THE LEFT SONS WHEN ON THE
C LEFT HALF OF THE TREE, AND TO STORE THE SIZES OF THE LEFT

```

```

C SONS WHILE TRAVERSING THE RIGHT SONS(DOWN THE RIGHT HALF)
  LOGICAL MAXE,MINE
  IF(M.EQ.1)GOTO 22
  IL(1)=2
  IL(2)=1
  KN=IN(1)
  KX=IN(2)
  IF(M.EQ.2)GOTO 21
  MP1=M+1
  MIN=1
  MX=1
  KX=IN(1)
  MAXE=.FALSE.
  MINE=.FALSE.
C FIND TWO VERTICES OF THE CONVEX HULL FOR THE INITIAL
C PARTITION
  DO 6 I=2,M
    J=IN(I)
    IF(X(1,J)-X(1,KX))3,1,2
1    MAXE=.TRUE.
    GOTO 3
2    MAXE=.FALSE.
    MX=I
    KX=J
3    IF(X(1,J)-X(1,KN))5,4,6
4    MINE=.TRUE.
    GOTO 6
5    MINE=.FALSE.
    MIN=I
    KN=J
6    CONTINUE
C IF THE MAX AND MIN ARE EQUAL, ALL M POINTS LIE ON A
C VERTICAL LINE
  IF(KX.EQ.KN)GOTO 18
C IF MAXE (OR MINE) HAS THE VALUE TRUE THERE ARE SEVERAL
C MAXIMA (OR MINIMA) WITH EQUAL FIRST COORDINATES
  IF(MAXE.OR.MINE)GOTO 23
7    IH(1)=KX
    IH(2)=KN
    NH=3
    INH=1
    NIB=1
    MA=M
    IN(MX)=IN(M)
    IN(M)=KX
    MM=M-2
    IF(MIN.EQ.M)MIN=MX
    IN(MIN)=IN(M-1)
    IN(M-1)=KN
C BEGIN BY PARTITIONING THE ROOT OF THE TREE
  CALL SPLIT(N,X,MM,IN,IH(1),IH(2),Ø,IA,MB,MXA,IB,IA(MA),
1  MXBB)
C FIRST TRAVERSE THE LEFT HALF OF THE TREE
C START WITH THE LEFT SON
8    NIB=NIB+IA(MA)
    MA=MA-1
9    IF(MXA.EQ.Ø)GOTO 11
    IL(NH)=IL(INH)
    IL(INH)=NH
    IH(NH)=IA(MXA)
    IA(MXA)=IA(MB)
    MB=MB-1
    NH=NH+1
    IF(MB.EQ.Ø)GOTO 1Ø
    ILINH=IL(INH)
    CALL SPLIT(N,X,MB,IA,IH(INH),IH(ILINH),1,IA,MBB,MXA,
1  IB(NIB),IA(MA),MXB)
    MB=MBB
    GOTO 8
C THEN THE RIGHT SON
1Ø   INH=IL(INH)
11   INH=IL(INH)
    MA=MA+1
    NIB=NIB-IA(MA)
    IF(MA.GE.M)GOTO 12
    IF(IA(MA).EQ.Ø)GOTO 11
    ILINH=IL(INH)

```



```

C ON THE LEFT SIDE OF THE TREE, THE RIGHT SON OF A RIGHT SON
C MUST REPRESENT A SUBSET OF POINTS WHICH IS INSIDE A
C TRIANGLE WITH VERTICES WHICH ARE ALSO VERTICES OF THE
C CONVEX POLYGON AND HENCE THE SUBSET MAY BE NEGLECTED.
  CALL SPLIT(N,X,IA(MA),IB(NIB),IH(INH),IH(ILINH),2,IA,
  1 MB,MXA,IB(NIB),MBB,MXB)
  IA(MA)=MBB
  GOTO 9
C NOW TRAVERSE THE RIGHT HALF OF THE TREE
12  MXB=MXBB
  MA=M
  MB=IA(MA)
  NIA=1
  IA(MA)=Ø
C START WITH THE RIGHT SON
13  NIA=NIA+IA(MA)
  MA=MA-1
14  IF(MXB.EQ.Ø)GOTO 16
  IL(NH)=IL(INH)
  IL(INH)=NH
  IH(NH)=IB(MXB)
  IB(MXB)=IB(MB)
  MB=MB-1
  NH=NH+1
  IF(MB.EQ.Ø)GOTO 15
  ILLINH=IL(INH)
  CALL SPLIT(N,X,MB,IB(NIB),IH(INH),IH(ILINH),-1,IA(NIA),
  1 IA(MA),MXA,IB(NIB),MBB,MXB)
  MB=MBB
  GOTO 13
C THEN THE LEFT SON
15  INH=IL(INH)
16  INH=IL(INH)
  MA=MA+1
  NIA=NIA-IA(MA)
  IF(MA.EQ.MP1)GOTO 17
  IF(IA(MA).EQ.Ø)GOTO 16
  ILLINH=IL(INH)
C ON THE RIGHT SIDE OF THE TREE, THE LEFT SON OF A LEFT SON
C MUST REPRESENT A SUBSET OF POINTS WHICH IS INSIDE A
C TRIANGLE WITH VERTICES WHICH ARE ALSO VERTICES OF THE
C CONVEX POLYGON AND HENCE THE SUBSET MAY BE NEGLECTED.
  CALL SPLIT(N,X,IA(MA),IA(NIA),IH(INH),IH(ILINH),-2,
  1 IA(NIA),MBB,MXA,IB(NIB),MB,MXB)
  GOTO 14
17  NH=NH-1
  RETURN
C ALL THE SPECIAL CASES ARE HANDLED DOWN HERE
C IF ALL THE POINTS LIE ON A VERTICAL LINE
18  KX=IN(1)
  KN=IN(1)
  DO 2Ø I=1,M
    J=IN(I)
    IF(X(2,J).LE.X(2,KX))GOTO 19
    MX=I
    KX=J
19  IF(X(2,J).GE.X(2,KN))GOTO 2Ø
  MIN=I
  KN=J
2Ø  CONTINUE
  IF(KX.EQ.KN)GOTO 22
C IF THERE ARE ONLY TWO POINTS
21  IH(1)=KX
  IH(2)=KN
  NH=3
  IF((X(1,KN).EQ.X(1,KX)).AND.(X(2,KN).EQ.X(2,KX)))NH=2
  GOTO 17
C IF THERE IS ONLY ONE POINT
22  NH=2
  IH(1)=IN(1)
  IL(1)=1
  GOTO 17
C MULTIPLE EXTREMES ARE HANDLED HERE
C IF THERE ARE SEVERAL POINTS WITH THE (SAME) LARGEST
C FIRST COORDINATE
23  IF(.NOT.MAXE)GOTO 25

```

```

      DO 24 I=1,M
        J=IN(I)
        IF(X(1,J).NE.X(1,KX))GOTO 24
        IF(X(2,J).LE.X(2,KX))GOTO 24
        MX=I
        KX=J
24     CONTINUE
C IF THERE ARE SEVERAL POINTS WITH THE (SAME) SMALLEST
C FIRST COORDINATE
25     IF(.NOT.MINE)GOTO 7
        DO 26 I=1,M
          J=IN(I)
          IF(X(1,J).NE.X(1,KN))GOTO 26
          IF(X(2,J).GE.X(2,KN))GOTO 26
          MIN=I
          KN=J
26     CONTINUE
      GOTO 7
      END

```

```

20
2.0      0.0
1.73     -1.0
1.0      1.73
0.0      2.0
0.1      0.1
-1.0     -1.73
0.2      -0.2
-1.73    1.0
-0.3     0.3
0.0      -2.0
-0.4     -0.4
-2.0     0.0
0.5      0.5
1.73     1.0
0.6      -0.6
-1.0     1.73
-0.7     0.7
-1.73    -1.0
-0.8     -0.8
1.0      -1.73

```

ALGORITHM 524

MP, A Fortran Multiple-Precision Arithmetic Package [A1]

RICHARD P. BRENT
Australian National University

Key Words and Phrases: arithmetic, multiple precision, extended precision, floating point, elementary function evaluation, Euler's constant, gamma function, polyalgorithm, software package, Fortran, machine-independent software, special function evaluation, Bessel functions, exponential integral, logarithmic intergral, Bernoulli numbers, zeta function, portable software
CR Categories: 3.15, 4.49, 5.11, 5.12, 5.15, 5.19, 5.25
Language: Fortran

DESCRIPTION

The design of the package and the theoretical background for the algorithms used are given in [1]. Details of calling sequences, etc., are given in the comments included here and in [2].

REFERENCES

1. BRENT, R.P. A Fortran multiple-precision arithmetic package. *ACM Trans. Math. Software* 4, 1 (March 1978), 57-70.
2. BRENT, R.P. MP users guide. Tech. Rep. 54, Computer Centre, Australian National U., Canberra, Australia, Sept. 1976.

ALGORITHM

[Only that portion of the listing which gives the introductory comments and a small example program is printed here. The complete listing, together with a Users' Guide giving further details, is available from the ACM Algorithms Distribution Service.

```

C ABCDEFGHIJKLMNOPQRSTUVWXYZ$0123456789+*/*( ), .
C
C $$          ***** COMMENTS *****
C
C DESCRIPTION OF MP (VERSION OF 17 FEBRUARY 1977)
C *****
C
C MP IS A MULTIPLE-PRECISION ARITHMETIC PACKAGE.
C IT IS ALMOST COMPLETELY MACHINE-INDEPENDENT, AND SHOULD

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```

MP000010
MP000020
MP000030
MP000040
MP000051
MP000060
MP000070
MP000080
MP000090

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```

C RUN ON ANY MACHINE WITH AN ANSI STANDARD FORTRAN COMPILER,
C SUFFICIENT MEMORY, AND A WORDLENGTH OF AT LEAST 16 BITS.
C SOME MODIFICATIONS WOULD BE NECESSARY FOR A WORDLENGTH
C OF LESS THAN 16 BITS.
C
C MP WAS WRITTEN BY R. BRENT (COMPUTER CENTRE, AUSTRALIAN
C NATIONAL UNIVERSITY), IN NOVEMBER 1973 (VERSION 1).
C CORRESPONDENCE SHOULD BE SENT TO R. P. BRENT, COMPUTER
C CENTRE, ANU, BOX 4, CANBERRA, ACT 2600, AUSTRALIA.
C
C MP HAS BEEN TESTED ON A UNIVAC 1108 (E LEVEL FORTRAN V),
C A UNIVAC 1100/42 (E AND T LEVEL FORTRAN V, ASCII FORTRAN,
C AND RALPH), A PDP 10 (FORTRAN 10(/NOOPT) AND FORTRAN 40),
C AN IBM 360/50 (FORTRAN G AND FORTRAN H, OPT = 2),
C AN IBM 360/91 AND 370/168 (FORTRAN H EXTENDED, OPT = 2),
C A CYBER 76 (FTN 4.2, OPT = 1) AND A PDP 11/45 (DOS).
C THESE MACHINES HAVE EFFECTIVE INTEGER WORDLENGTHS RANGING
C FROM 16 TO 48 BITS.
C
C MP WORKS WITH NORMALIZED FLOATING-POINT NUMBERS.
C THE BASE (B) AND NUMBER OF DIGITS (T) ARE
C ARBITRARY (SUBJECT TO SOME RESTRICTIONS GIVEN BELOW),
C AND MAY BE VARIED DYNAMICALLY.
C
C T-DIGIT FLOATING-POINT NUMBERS ARE STORED IN INTEGER ARRAYS OF
C DIMENSION T+2, WITH THE FOLLOWING CONVENTIONS -
C   WORD 1 = SIGN (0, -1 OR +1)
C   WORD 2 = EXPONENT (TO BASE B)
C   WORDS 3 TO T+2 = NORMALIZED FRACTION
C NOTE THAT WORDS 2 TO T+2 ARE UNDEFINED IF SIGN = 0.
C
C ARITHMETIC IS ROUNDED, AND FOUR GUARD DIGITS ARE USED
C FOR ADDITION AND MULTIPLICATION, SO THE CORRECTLY ROUNDED
C RESULT IS USUALLY PRODUCED. DIVISION, SQRT ETC ARE DONE
C BY NEWTONS METHOD, BUT GIVE THE EXACT RESULT IF IT CAN BE
C REPRESENTED WITH T-2 DIGITS. OTHER ROUTINES (MPSIN, MPLN ETC)
C USUALLY GIVE A RESULT  $Y = F(X)$  WHICH COULD BE OBTAINED
C BY MAKING AN  $O(B^{*(1-T)})$  PERTURBATION IN X, EVALUATING F
C EXACTLY, THEN MAKING AN  $O(B^{*(1-T)})$  PERTURBATION IN Y.
C
C EXPONENTS CAN LIE IN THE RANGE -M, ... , +M INCLUSIVE,
C WHERE M IS SET BY THE USER.
C ON UNDERFLOW DURING AN ARITHMETIC OPERATION, THE RESULT
C IS SET TO ZERO BY SUBROUTINE MPUNFL.
C ON OVERFLOW SUBROUTINE MPOVFL IS CALLED AND EXECUTION
C IS TERMINATED WITH AN ERROR MESSAGE.
C ERROR MESSAGES ARE PRINTED ON LOGICAL UNIT LUN, WHERE LUN
C IS SET BY THE USER, AND THEN EXECUTION IS TERMINATED
C BY A CALL TO SUBROUTINE MPERR. IT IS ASSUMED THAT LOGICAL
C RECORDS OF UP TO 80 CHARACTERS MAY BE WRITTEN ON UNIT LUN.
C A WORKING ARRAY OF SIZE MXR (SEE BELOW) MUST BE PROVIDED
C IN COMMON.
C
C THE PARAMETERS B, T, M, LUN AND MXR ARE PASSED TO THE UTILITY
C ROUTINES IN COMMON, TOGETHER WITH A WORKING ARRAY R WHICH
C MUST BE SUFFICIENTLY LARGE (SEE BELOW). MOST ROUTINES
C USE THE STATEMENTS -
C   COMMON B, T, M, LUN, MXR, R
C   INTEGER B, T, R(1)
C AND IT IS ASSUMED THAT R IS DIMENSIONED SUFFICIENTLY LARGE
C IN THE CALLING PROGRAM, AND THAT MXR IS SET TO THE
C DIMENSION OF R IN THE CALLING PROGRAM.
C WARNING - IT IS ASSUMED THAT THE COMPILER PASSES ADDRESSES OF
C ***** ARRAYS USED AS ARGUMENTS IN SUBROUTINE CALLS (I.E., CALL
C BY REFERENCE), AND DOES NOT CHECK FOR ARRAY BOUNDS
C VIOLATIONS (EITHER FOR ARGUMENTS OR FOR ARRAYS IN COMMON).
C APART FROM THESE VIOLATIONS, MP IS WRITTEN ENTIRELY IN
C ANSI STANDARD FORTRAN (ANSI X3.9-1966). THIS HAS BEEN
C CHECKED BY THE PFORT VERIFIER.
C
C RESTRICTIONS -
C   B (THE BASE) MUST BE AT LEAST 2,
C   T (NUMBER OF DIGITS) MUST BE AT LEAST 2,
C   M (EXPONENT RANGE) MUST BE GREATER THAN T AND LESS
C   THAN 1/4 THE LARGEST MACHINE-REPRESENTABLE INTEGER,
C    $8*B**2-1$  MUST BE NO GREATER THAN THE LARGEST MACHINE-

```

```

MP000100
MP000110
MP000120
MP000130
MP000140
MP000150
MP000160
MP000170
MP000180
MP000190
MP000200
MP000210
MP000220
MP000230
MP000240
MP000250
MP000260
MP000270
MP000280
MP000290
MP000300
MP000310
MP000320
MP000330
MP000340
MP000350
MP000360
MP000370
MP000380
MP000390
MP000400
MP000410
MP000420
MP000430
MP000440
MP000450
MP000460
MP000470
MP000480
MP000490
MP000500
MP000510
MP000520
MP000530
MP000540
MP000550
MP000560
MP000570
MP000580
MP000590
MP000600
MP000610
MP000620
MP000630
MP000640
MP000650
MP000660
MP000670
MP000680
MP000690
MP000700
MP000710
MP000720
MP000730
MP000740
MP000750
MP000760
MP000770
MP000780
MP000790
MP000800
MP000810
MP000820
MP000830
MP000840
MP000850

```

C	REPRESENTABLE INTEGER, AND THE INTEGERS 0, 1, ... , B	MP000860
C	MUST BE EXACTLY REPRESENTABLE AS SINGLE-PRECISION FLOATING-	MP000870
C	POINT NUMBERS, AND	MP000880
C	B**(T-1) SHOULD BE AT LEAST 10**7.	MP000890
C		MP000900
C	B AND T MAY BE SET TO GIVE THE EQUIVALENT OF A SPECIFIED	MP000910
C	NUMBER OF DECIMAL PLACES BY CALLING MPSET (SEE BELOW), OR MAY	MP000920
C	BE SET DIRECTLY BY THE USER. IF MPSET IS NOT CALLED, THE USER	MP000930
C	MUST REMEMBER TO INITIALIZE M, LUN AND MXR (SEE ABOVE) AS WELL	MP000940
C	AS B AND T BEFORE CALLING ANY MP ROUTINES.	MP000950
C		MP000960
C	FOR EFFICIENCY CHOOSE B FAIRLY LARGE, SUBJECT TO THE RESTRICTIONS	MP000970
C	GIVEN ABOVE. FOR EXAMPLE, IF THE WORDLENGTH IS	MP000980
C	48 BITS, COULD USE B = 4194304 OR 1000000,	MP000990
C	36 BITS, COULD USE B = 65536 OR 10000,	MP001000
C	32 BITS, COULD USE B = 16384 OR 10000,	MP001010
C	24 BITS, COULD USE B = 1024 OR 1000,	MP001020
C	18 BITS, COULD USE B = 128 OR 100,	MP001030
C	16 BITS, COULD USE B = 64 OR 10.	MP001040
C		MP001050
C	AVOID MULTIPLICATION OR DIVISION BY MP NUMBERS, AS	MP001060
C	THESE TAKE O(T**2) OPERATIONS, WHEREAS MULT./DIV. BY	MP001070
C	INTEGERS TAKE O(T) OPERATIONS.	MP001080
C		MP001090
C	MP NUMBERS USED AS ARGUMENTS OF SUBROUTINES NEED NOT BE	MP001100
C	DISTINCT. FOR EXAMPLE,	MP001110
C	CALL MPADD (X, Y, Y) OR CALL MPEXP (X, X) ARE OK.	MP001120
C	HOWEVER, DISTINCT ARRAYS WHICH OVERLAP SHOULD NOT BE USED.	MP001130
C		MP001140
C	FOR ADDITIONAL DETAILS SEE - A FORTRAN MULTIPLE-PRECISION	MP001150
C	ARITHMETIC PACKAGE (BY R. P. BRENT), TO APPEAR IN ACM	MP001160
C	TRANSACTIONS ON MATHEMATICAL SOFTWARE (AVAILABLE AS A	MP001170
C	CARNEGIE-MELLON UNIV. COMPUTER SCIENCE DEPT. REPORT,	MP001180
C	PITTSBURGH, PENNSYLVANIA, MAY 1976) AND THE MP USERS GUIDE.	MP001190
C		MP001200
C	SUMMARY OF MP ROUTINES	MP001210
C	*****	MP001220
C		MP001230
C	BASIC ARITHMETIC - MPADD, MPADDI, MPADDQ, MPDIV, MPDIVI,	MP001240
C	MPMUL, MPMULI, MPMULQ, MPREC, MPSUB	MP001250
C		MP001260
C	POWERS AND ROOTS - MPPWR, MPPWR2, MPQPWR, MPROOT, MPSQRT	MP001270
C		MP001280
C	ELEMENTARY FUNCTIONS - MPASIN, MPATAN, MPCOS, MPCOSH, MPEXP,	MP001290
C	MPLN, MPLNGS, MPLNI, MPSIN, MPSINH, MPTAN,	MP001300
C	MPTANH	MP001310
C		MP001320
C	SPECIAL FUNCTIONS - MPBESJ, MPDAW, MPEI, MPERF, MPERFC,	MP001330
C	MPGAM, MPGAMQ, MPLI, MPLNGM	MP001340
C		MP001350
C	CONSTANTS - MPBERN, MPEPS, MPEUL, MPMAXR, MPMINR, MPPI,	MP001360
C	MPPIGL, MPZETA	MP001370
C		MP001380
C	INPUT AND OUTPUT - MPDUMP, MPIN, MPINE, MPOUT, MPOUTE, MPOUT2	MP001390
C		MP001400
C	CONVERSION - MPCDM, MPCIM, MPCMD, CPCMDE, MPCMEF, MPCMI,	MP001410
C	MPCMIM, MPCMR, MPCMRE, MPCQM, MPCRM	MP001420
C		MP001430
C	COMPARISON - MPCMPA, MPCMPI, MPCMPR, MPCOMP	MP001440
C		MP001450
C	GENERAL UTILITY ROUTINES - MPABS, MPCLR, MPCMF, MPMAX, MPMIN,	MP001460
C	MPNEG, MPPACK, MPPOLY, MPSET, MPSTR, MPUNPK	MP001470
C		MP001480
C	ERROR DETECTION AND HANDLING - MPCHK, MPERR, MPOVFL, MPUNFL	MP001490
C		MP001500
C	TEST PROGRAMS - EXAMPLE, TEST, TESTV, TEST2	MP001510
C		MP001520
C	MISCELLANEOUS ROUTINES USED BY THE ABOVE - MPADD2, MPADD3,	MP001530
C	MPART1, MPBES2, MPERF2, MPERF3, MPEXP1,	MP001540
C	MPEXT, MPGCD, MPHANK, MPLNS, MPL235,	MP001550
C	MPMLP, MPMUL2, MPNZR, MPSIN1, MP40D,	MP001560
C	MP40E, MP40F, MP40G, TIMEMP	MP001570
C		MP001580
C	LIST OF MP ROUTINES	MP001590
C	*****	MP001600
C		MP001610

C THE ROUTINES PROVIDED IN MP ARE LISTED BELOW. FOR MORE DETAILS OF
C CALLING SEQUENCES, RESTRICTIONS, ACCURACY, AND ERROR CONDITIONS,
C SEE THE COMMENTS IN EACH ROUTINE. SPACE REQUIRED (I.E. DIMENSION
C OF R IN COMMON) IS T+4 WORDS UNLESS NOTED BELOW. THE ROUTINES
C INDICATED BY * ARE THOSE MOST LIKELY TO BE OF INTEREST TO THE USER.
C FOR FAST EXECUTION THE ROUTINES WHICH SHOULD BE OPTIMIZED ARE
C MPNZR, MPMLP, MPDIVI, MPADD2, MPADD3, AND MPMUL2.

C
C IN THE LIST BELOW AN MP NUMBER MEANS A MULTIPLE-PRECISION
C NUMBER AS DESCRIBED ABOVE, M(T) MEANS THE TIME TO MULTIPLY TWO
C T-DIGIT MP NUMBERS (SEE MPMUL), AN INTEGER MEANS A SINGLE-PRECISION
C INTEGER, A RATIONAL NUMBER MEANS THE RATIO OF TWO INTEGERS.
C X, Y, ... ARE MP NUMBERS, I, J, ... ARE INTEGERS,
C AND RI, RX, ... ARE SINGLE-PRECISION REAL NUMBERS.
C TIME BOUNDS SUCH AS O(T**2) ARE AS T TENDS TO INFINITY
C WITH EVERYTHING ELSE FIXED.

C
C * EXAMPLE A SMALL MAIN PROGRAM GIVING AN EXAMPLE OF THE USE OF MP.
C * MPABS COMPUTES ABSOLUTE VALUE OF AN MP NUMBER
C CALL MPABS (X, Y) MEANS Y = ABS(X)
C * MPADD ADDS TWO MP NUMBERS
C CALL MPADD (X, Y, Z) MEANS Z = X + Y
C * MPADDI ADDS AN MP NUMBER TO AN INTEGER,
C GIVING A MULTIPLE-PRECISION RESULT, SPACE = 2T+6
C CALL MPADDI (X, IY, Z) MEANS Z = X + IY
C * MPADDQ ADDS A RATIONAL NUMBER TO AN MP NUMBER,
C SPACE = 2T+6
C CALL MPADDQ (X, I, J, Y) MEANS Y = X + I/J
C MPADD2 ROUTINE CALLED BY MPADD AND MPSUB
C MPADD3 ROUTINE CALLED BY MPADD2
C MPART1 COMPUTES ARCTAN(1/N) FOR N .GT. 1 (CALLED BY MPPI)
C SPACE = 2T+6
C * MPASIN COMPUTES ARCSIN OF AN MP NUMBER,
C USING AN O(M(T)T) METHOD, SPACE = 5T+12
C CALL MPASIN (X, Y) MEANS Y = ARCSIN(X)
C * MPATAN COMPUTES ARCTAN OF AN MP NUMBER
C USING AN O(T.M(T)) ALGORITHM, SPACE = 5T+12
C CALL MPATAN (X, Y) MEANS Y = ARCTAN(X)
C * MPBERN COMPUTES BERNOULLI NUMBERS B2, B4, B6, ...
C SPACE = 8T+18
C * MPBESJ COMPUTES BESSEL FUNCTION J(NU,X) FOR MP X
C AND SMALL INTEGER NU, SPACE = 14T+156
C CALL MPBESJ (X, NU, Y) MEANS Y = J(NU,X)
C MPBES2 ROUTINE CALLED BY MPBESJ (USES BACKWARD RECURRENCE
C TO EVALUATE J(NU,X)), SPACE = 8T+18
C MPCDM CONVERTS DOUBLE-PRECISION TO MULTIPLE-PRECISION
C MPCHK PRINTS ERROR MESSAGE ON UNIT LUN IF B, T, M OR MXR
C IS ILLEGAL, OR ON UNIT 6 IF LUN IS ILLEGAL
C (LUN SHOULD BE IN RANGE 1 TO 99)
C * MPCIM CONVERTS INTEGER TO MULTIPLE-PRECISION
C CALL MPCIM (IX, Z) MEANS Z = IX
C MPCLR SETS SOME DIGITS OF AN MP NUMBER TO ZERO
C MPCMD CONVERTS AN MP NUMBER TO DOUBLE-PRECISION REAL
C MPCMDE CONVERTS AN MP NUMBER TO (DOUBLE-PRECISION)
C FRACTION AND (DECIMAL) EXPONENT,
C SPACE = 6T+14
C MPCMEF CONVERTS MP NUMBER TO FRACTION AND (DECIMAL)
C EXPONENT, SPACE = 5T+12
C * MPCMF FINDS FRACTIONAL PART OF AN MP NUMBER
C MPCMI CONVERTS AN MP NUMBER TO AN INTEGER
C * MPCMIM CONVERTS AN MP NUMBER TO A MULTIPLE-PRECISION INTEGER
C * MPCMPA COMPARES ABSOLUTE VALUES OF TWO MP NUMBERS
C MPCMPA (X, Y) RETURNS SIGN(ABS(X)-ABS(Y))
C * MPCMPI COMPARES AN MP NUMBER WITH AN INTEGER, SPACE = 2T+6
C MPCMPI (X, I) RETURNS SIGN(X-I)
C * MPCMPR COMPARES AN MP NUMBER WITH A REAL, SPACE = 2T+6
C MPCMPR (X, RI) RETURNS SIGN(X-RI)
C * MPCMR CONVERTS AN MP NUMBER TO (SINGLE-PRECISION) REAL
C CALL MPCMR (X, RZ) MEANS RZ = SNGL(X)
C * MPCMRE CONVERTS AN MP NUMBER TO EXPONENT AND
C (SINGLE-PRECISION) FRACTION, I.E. F*10**I
C SPACE = 6T+14
C * MPCOMP COMPARES TWO MP NUMBERS
C MPCOMP (X, Y) RETURNS SIGN(X-Y)
C * MPCOS COMPUTES COSINE OF AN MP NUMBER, USING AN
C O(M(T)T/LOG(T)) METHOD, SPACE = 5T+12

MP001620
MP001630
MP001640
MP001650
MP001660
MP001670
MP001680
MP001690
MP001700
MP001710
MP001720
MP001730
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MP001890
MP001900
MP001910
MP001920
MP001930
MP001940
MP001950
MP001960
MP001970
MP001980
MP001990
MP002000
MP002010
MP002020
MP002030
MP002040
MP002050
MP002060
MP002070
MP002080
MP002090
MP002100
MP002110
MP002120
MP002130
MP002140
MP002150
MP002160
MP002170
MP002180
MP002190
MP002200
MP002210
MP002220
MP002230
MP002240
MP002250
MP002260
MP002270
MP002280
MP002290
MP002300
MP002310
MP002320
MP002330
MP002340
MP002350
MP002360
MP002370

C		CALL MPCOS (X, Y) MEANS $Y = \cos(X)$	MP002380
C *	MPCOSH	COMPUTES HYPERBOLIC COSINE OF AN MP NUMBER	MP002390
C		USING MPEXP, SPACE = 5T+12	MP002400
C		CALL MPCOSH (X, Y) MEANS $Y = \cosh(X)$	MP002410
C *	MPCQM	CONVERTS A RATIONAL NUMBER TO MULTIPLE-PRECISION	MP002420
C		CALL MPCQM (I, J, Q) MEANS $Q = I/J$	MP002430
C	MPCRM	CONVERTS REAL TO MULTIPLE-PRECISION	MP002440
C		CALL MPCRM (RX, Z) MEANS $Z = RX$	MP002450
C *	MPDAW	COMPUTES DAWSONS INTEGRAL, $DAW(X) = \exp(-X^2) * (\text{INTEGRAL FROM } 0 \text{ TO } X \text{ OF } \exp(U^2) \text{ DU})$, SPACE = 5T+17	MP002460
C		CALL MPDAW (X, Y) MEANS $Y = DAW(X)$	MP002480
C *	MPDIV	DIVIDES TWO MP NUMBERS, SPACE = 4T+10	MP002490
C		CALL MPDIV (X, Y, Z) MEANS $Z = X/Y$	MP002500
C *	MPDIVI	DIVIDES AN MP NUMBER BY AN INTEGER	MP002510
C		USING AN O(T) METHOD (MUCH FASTER THAN MPDIV)	MP002520
C		CALL MPDIVI (X, IY, Z) MEANS $Z = X/IY$	MP002530
C	MPDUMP	DUMPS AN MP NUMBER (USEFUL FOR DEBUGGING)	MP002540
C		CALL MPDUMP (X) DUMPS THE MP NUMBER X ON UNIT LUN	MP002550
C *	MPEI	EVALUATES EXPONENTIAL INTEGRAL OF AN MP NUMBER, SPACE = 19T+31	MP002560
C		CALL MPEI (X, Y) MEANS $Y = Ei(X)$	MP002570
C *	MPEPS	COMPUTES THE (MULTIPLE-PRECISION) MACHINE PRECISION	MP002580
C		CALL MPEPS (X) MEANS $X = 0.5 * B^{(1-T)}$ IF B EVEN	MP002600
C *	MPERF	COMPUTES ERROR FUNCTION OF AN MP NUMBER, SPACE = 5T+12	MP002610
C		CALL MPERF (X, Y) MEANS $Y = \text{ERF}(X)$	MP002620
C *	MPERFC	COMPUTES COMPLEMENTARY ERROR FUNCTION OF AN MP NUMBER, SPACE = 12T+26	MP002630
C		CALL MPERFC (X, Y) MEANS $Y = \text{ERFC}(X)$	MP002640
C	MPERF2	COMPUTES $\exp(X^2) * (\text{INTEGRAL FROM } 0 \text{ TO } X \text{ OF } \exp(-U^2) \text{ DU})$, CALLED BY MPERF, SPACE = 5T+12	MP002650
C		ROUTINE CALLED BY MPERF, MPDAW AND MPERFC, SPACE = 4T+10	MP002660
C	MPERF3		MP002670
C			MP002680
C			MP002690
C			MP002700
C	MPERR	ERROR HANDLING ROUTINE (TERMINATES EXECUTION AT PRESENT BUT MAY EASILY BE MODIFIED).	MP002710
C			MP002720
C *	MPEUL	RETURNS EULERS CONSTANT ($\text{GAMMA} = 0.57721566\dots$) TO MULTIPLE-PRECISION ACCURACY, SPACE = 5T+14	MP002730
C		CALL MPEUL (G) MEANS $G = 0.57721566\dots$	MP002740
C *	MPEXP	COMPUTES EXPONENTIAL OF A MULTIPLE-PRECISION NUMBER, USING AN $O(\sqrt{T}M(T))$ METHOD, SPACE = 4T+10	MP002750
C		CALL MPEXP (X, Y) MEANS $Y = \exp(X)$	MP002760
C	MPEXP1	COMPUTES $\exp(X) - 1$ FOR $\text{ABS}(X) \leq 1$ (CALLED BY MPEXP, MPSINH AND MPTANH), SPACE = 3T+8	MP002770
C			MP002780
C	MPEXT	A ROUNDING ROUTINE CALLED BY MPDIV AND MPSQRT	MP002790
C *	MPGAM	COMPUTES GAMMA FUNCTION OF AN MP ARGUMENT, SPACE SAME AS FOR MPLNGM (IN WORST CASE)	MP002800
C		CALL MPGAM (X, Y) MEANS $Y = \text{GAMMA}(X)$	MP002810
C	MPGAMQ	COMPUTES GAMMA FUNCTION OF A RATIONAL ARGUMENT, USING AN $O(T^2)$ METHOD, SPACE = 6T+12	MP002820
C		CALL MPGAMQ (I, J, X) MEANS $X = \text{GAMMA}(I/J)$	MP002830
C	MPGCD	DIVIDES TWO INTEGERS BY THEIR GREATEST COMMON DIVISOR (CALLED BY MPMULQ, MPGAMQ, ETC)	MP002840
C			MP002850
C	MPHANK	ROUTINE CALLED BY MPBESJ (EVALUATES HANKELS ASYMPTOTIC SERIES FOR BESSEL FUNCTIONS), SPACE = 11T+24	MP002860
C			MP002870
C *	MPIN	CONVERTS FIXED-POINT NUMBER READ UNDER A1 FORMAT TO MULTIPLE-PRECISION, SPACE = 3T+11	MP002880
C			MP002890
C *	MPINE	SAME AS MPIN BUT RESULT IS MULTIPLIED BY A POWER OF TEN (USEFUL FOR READING IN FLOATING-POINT NUMBERS), SPACE = 5T+12	MP002900
C			MP002910
C *	MPLI	EVALUATES LOGARITHMIC INTEGRAL $Li(X)$, SPACE = 19T+31	MP002920
C		CALL MPLI (X, Y) MEANS $Y = Li(X)$	MP002930
C *	MPLN	COMPUTES NATURAL LOG OF AN MP NUMBER, USING AN $O(\sqrt{T}M(T))$ METHOD, SPACE = 6T+14	MP002940
C		CALL MPLN (X, Y) MEANS $Y = \ln(X)$	MP002950
C	MPLNGM	COMPUTES $\ln(\text{GAMMA}(X))$ FOR POSITIVE MP X, USING STIRLINGS APPROXIMATION, SPACE = $11T + 24 + NL * ((T+3)/2)$, WHERE NL IS THE NUMBER OF TERMS USED IN THE ASYMPTOTIC EXPANSION, NL .LE. $(2 + T * \ln(B)/8)$	MP002960
C			MP002970
C		CALL MPLNGM (X, Y) MEANS $Y = \ln(\text{GAMMA}(X))$	MP002980
C	MPLNGS	COMPUTES NATURAL LOG OF AN MP NUMBER, USING THE GAUSS-SALAMIN ALGORITHM. RECOMMENDED FOR TESTING MPLN AND MPLNI ONLY (UNLESS T LARGE). SPACE = 6T+26	MP002990
C			MP003000
C *	MPLNI	COMPUTES NATURAL LOG OF AN INTEGER, USING AN $O(T^2)$ METHOD (FASTER THAN MPLN), SPACE = 3T+8	MP003010
C			MP003020
			MP003030
			MP003040
			MP003050
			MP003060
			MP003070
			MP003080
			MP003090
			MP003100
			MP003110
			MP003120
			MP003130

C		CALL MPLNI (N, X) MEANS $X = \text{LN}(N)$	MP003140
C	MPLNS	COMPUTES $\text{LN}(1+X)$ FOR SMALL MP X. SPACE = 5T+12	MP003150
C	MPL235	COMPUTES NATURAL LOG OF AN INTEGER WHOSE PRIME	MP003160
C		FACTORS ARE 2, 3 AND/OR 5 (CALLED BY MPLNI),	MP003170
C		SPACE = 3T+8	MP003180
C	* MPMAX	COMPUTES THE MAXIMUM OF TWO MP NUMBERS	MP003190
C		CALL MPMAX (X, Y, Z) MEANS $Z = \text{MAX}(X, Y)$	MP003200
C	* MPMAXR	COMPUTES THE LARGEST POSITIVE MP NUMBER	MP003210
C		CALL MPMAXR (X) MEANS X = MP NUMBER WITH EXPONENT M	MP003220
C		AND ALL DIGITS B-1	MP003230
C	* MPMIN	COMPUTES THE MINIMUM OF TWO MP NUMBERS	MP003240
C		CALL MPMIN (X, Y, Z) MEANS $Z = \text{MIN}(X, Y)$	MP003250
C	* MPMINR	RETURNS THE SMALLEST NORMALIZED POSITIVE MP NUMBER	MP003260
C		CALL MPMINR (X) MEANS $X = B^{*(-M-1)}$	MP003270
C	MPMLP	INNER LOOP ROUTINE CALLED BY MPMUL	MP003280
C	* MPMUL	MULTIPLIES TWO MP NUMBERS	MP003290
C		USING AN $M(T) = O(T^{**2})$ ALGORITHM	MP003300
C		CALL MPMUL (X, Y, Z) MEANS $Z = X*Y$	MP003310
C	* MPMULI	MULTIPLIES AN MP NUMBER BY AN	MP003320
C		INTEGER USING AN $O(T)$ METHOD (FASTER THAN MPMUL)	MP003330
C		CALL MPMULI (X, IY, Z) MEANS $Z = X*IY$	MP003340
C	* MPMULQ	MULTIPLIES MP NUMBER BY A RATIONAL NUMBER	MP003350
C		CALL MPMULQ (X, I, J, Y) MEANS $Y = X*I/J$	MP003360
C	MPMUL2	ROUTINE CALLED BY MPMULI	MP003370
C	* MPNEG	REVERSES SIGN OF AN MP NUMBER	MP003380
C		CALL MPNEG (X, Y) MEANS $Y = -X$	MP003390
C	MPNZR	NORMALIZES AND ROUNDS OR TRUNCATES (CALLED BY	MP003400
C		MPADD2, MPDIVI, MPMUL AND MPMUL2)	MP003410
C	* MPOUT	CONVERTS MULTIPLE-PRECISION TO A FORM SUITABLE FOR	MP003420
C		PRINTING UNDER A1 FORMAT (CORRESPONDS TO F OR I	MP003430
C		FORMATS), SPACE = 3T+11	MP003440
C	* MPOUTE	SIMILAR TO MPOUT BUT GIVES (DECIMAL) EXPONENT AND	MP003450
C		FRACTION (CORRESPONDS TO E FORMAT), SPACE = 6T+14	MP003460
C	MPOUT2	SAME AS MPOUT BUT ANY BASE FROM 2 TO 16 MAY BE	MP003470
C		USED FOR OUTPUT REPRESENTATION, SPACE = 3T+11	MP003480
C	MPOVFL	ROUTINE CALLED ON MULTIPLE-PRECISION OVERFLOW	MP003490
C		(CALLS MPERR AT PRESENT BUT EASILY MODIFIED)	MP003500
C	* MPPACK	PACKS MP NUMBERS INTO ARRAYS OF DIMENSION	MP003510
C		$(T+3)/2$ (USEFUL TO SAVE SPACE),	MP003520
C		UNPACKING MAY BE PERFORMED WITH MPUNPK	MP003530
C	* MPPI	RETURNS PI TO MULTIPLE-PRECISION ACCURACY,	MP003540
C		USING AN $O(T^{**2})$ METHOD, SPACE = 3T+8	MP003550
C		CALL MPPI (X) MEANS $X = 3.14159265\dots$	MP003560
C	MPPIGL	RETURNS PI TO MULTIPLE-PRECISION ACCURACY,	MP003570
C		USING GAUSS-LEGENDRE $O(\text{LOG}(T)M(T))$ METHOD,	MP003580
C		RECOMMENDED FOR TESTING MPPI ONLY, SPACE = 6T+14	MP003590
C	* MPPOLY	EVALUATES A POLYNOMIAL WITH INTEGER COEFFICIENTS,	MP003600
C		SPACE = 3T+8	MP003610
C	* MPPWR	RAISES MP NUMBER TO INTEGER POWER,	MP003620
C		SPACE = 4T+10	MP003630
C		CALL MPPWR (X, N, Y) MEANS $Y = X**N$	MP003640
C	* MPPWR2	RAISES NONNEGATIVE MP NUMBER TO MP POWER,	MP003650
C		SPACE = 7T+16	MP003660
C		CALL MPPWR2 (X, Y, Z) MEANS $Z = X**Y$	MP003670
C	* MPQPWR	RAISES RATIONAL NUMBER TO RATIONAL POWER,	MP003680
C		SPACE = 4T+10	MP003690
C		CALL MPQPWR (I, J, K, L, X) MEANS $X = (I/J)**(K/L)$	MP003700
C	* MPREC	FORMS RECIPROCAL OF MP NUMBER,	MP003710
C		USING NEWTONS METHOD, SPACE = 4T+10	MP003720
C		CALL MPREC (X, Y) MEANS $Y = 1/X$	MP003730
C	* MPROOT	COMPUTES THE N-TH ROOT OF AN MP NUMBER	MP003740
C		USING NEWTONS METHOD, SPACE = 4T+10	MP003750
C		CALL MPROOT (X, N, Y) MEANS $Y = X**(1/N)$	MP003760
C	* MPSET	SETS THE BASE B AND DIGITS T ETC GIVEN THE	MP003770
C		EQUIVALENT NUMBER OF DECIMAL PLACES REQUIRED	MP003780
C		WARNING - MAY CAUSE AN INTEGER OVERFLOW,	MP003790
C		***** FOR DETAILS SEE COMMENTS IN MPSET	MP003800
C	* MPSIN	COMPUTES SINE OF AN MP NUMBER,	MP003810
C		USING AN $O(M(T)T/\text{LOG}(T))$ METHOD, SPACE = 5T+12	MP003820
C		CALL MPSIN (X, Y) MEANS $Y = \text{SIN}(X)$	MP003830
C	* MPSINH	COMPUTES HYPERBOLIC SINE OF AN MP NUMBER,	MP003840
C		USING MPEXP, SPACE = 5T+12	MP003850
C		CALL MPSINH (X, Y) MEANS $Y = \text{SINH}(X)$	MP003860
C	MPSINI	COMPUTES $\text{SIN}(X)$ OR $\text{COS}(X)$ FOR $\text{ABS}(X) \leq 1$, CALLED	MP003870
C		BY MPSIN, MPCOS AND MPTAN, SPACE = 3T+8	MP003880
C	* MPSQRT	COMPUTES SQUARE ROOT OF A NONNEGATIVE MP NUMBER,	MP003890

C	USING NEWTONS METHOD, SPACE = 4T+10	MP003900	
C	CALL MPSQRT (X, Y) MEANS $Y = \sqrt{X}$	MP003910	
C * MPSTR	STORES ONE MP NUMBER IN ANOTHER	MP003920	
C	CALL MPSTR (X, Y) MEANS $Y = X$	MP003930	
C * MPSUB	SUBTRACTS ONE MP NUMBER FROM ANOTHER	MP003940	
C	CALL MPSUB (X, Y, Z) MEANS $Z = X - Y$	MP003950	
C * MPTAN	COMPUTES TAN OF AN MP NUMBER,	MP003960	
C	USING MPSINI, SPACE = 6T+20	MP003970	
C	CALL MPTAN (X, Y) MEANS $Y = \tan(X)$	MP003980	
C * MPTANH	COMPUTES HYPERBOLIC TAN OF AN MP NUMBER,	MP003990	
C	USING MPEXP, SPACE = 5T+12	MP004000	
C	CALL MPTANH (X, Y) MEANS $Y = \tanh(X)$	MP004010	
C	ROUTINE CALLED ON MULTIPLE-PRECISION UNDERFLOW	MP004020	
C	(SETS RESULT TO ZERO AT PRESENT BUT EASILY MODIFIED)	MP004030	
C * MPUNPK	UNPACKS AN ARRAY FORMED BY MPPACK TO GIVE AN MP	MP004040	
C	NUMBER IN STANDARD FORMAT	MP004050	
C * MPZETA	COMPUTES RIEMANN ZETA FUNCTION FOR POSITIVE	MP004060	
C	INTEGER ARGUMENTS	MP004070	
C	SPACE = 8T+18+NL*((T+3)/2), WHERE NL IS THE	MP004080	
C	NUMBER OF TERMS USED IN THE ASYMPTOTIC	MP004090	
C	EXPANSION, NL .LE. (1 + 0.1*T*LN(B))	MP004100	
C	CALL MPZETA (N, X) MEANS $X = \zeta(N)$	MP004110	
C * MP40D	OUTPUT ROUTINE CALLED BY TEST PROGRAM,	MP004120	
C	USEFUL FOR EASY FIXED-POINT OUTPUT,	MP004130	
C	SPACE = 3T+N+14 FOR N DECIMAL PLACE OUTPUT	MP004140	
C	CALL MP40D (N, X) WRITES X TO N DECIMAL PLACES ON UNIT	MP004150	
C	LUN, ASSUMING ABS(X) .LT. 10	MP004160	
C	MP40E	OUTPUT ROUTINE CALLED BY MP40D	MP004170
C * MP40F	OUTPUT ROUTINE CALLED BY TEST2 PROGRAM,	MP004180	
C	USEFUL FOR EASY FLOATING-POINT OUTPUT,	MP004190	
C	SPACE = 6T+N+17 FOR N SIGNIFICANT FIGURE OUTPUT	MP004200	
C	CALL MP40F (N, X) WRITES X TO N SIGNIFICANT FIGURES	MP004210	
C	(IN DECIMAL EXPONENT AND FRACTION FORM) ON UNIT LUN	MP004220	
C	MP40G	OUTPUT ROUTINE CALLED BY MP40F	MP004230
C * TEST	A MAIN PROGRAM WHICH TESTS SOME OF THE ROUTINES IN MP	MP004240	
C	WHILE COMPUTING VARIOUS CONSTANTS TO 40 DECIMAL PLACES	MP004250	
C	TESTV	A VERSION OF TEST WITH VARIABLE-PRECISION COMPUTATION	MP004260
C	AND OUTPUT	MP004270	
C * TEST2	ANOTHER TEST PROGRAM WHICH TESTS ROUTINES	MP004280	
C	NOT CALLED BY TEST OR TESTV	MP004290	
C	TIMEMP	A MACHINE-DEPENDENT FUNCTION CALLED BY TESTV, SHOULD	MP004300
C	BE MODIFIED BY THE USER BEFORE TESTV IS RUN.	MP004310	
C		MP004320	
C	INDEX	MP004330	
C	*****	MP004340	
C		MP004350	
C	THE STARTING LINE SEQUENCE NUMBERS (GIVEN IN COLUMNS 73-80)	MP004360	
C	OF THE MP ROUTINES ARE AS FOLLOWS.	MP004370	
C		MP004380	
C	COMMENTS	MP000030	
C	EXAMPLE	MP000540	
C	MPABS	MP000610	
C	MPADD	MP000620	
C	MPADDI	MP000630	
C	MPADDQ	MP000650	
C	MPADD2	MP000660	
C	MPADD3	MP000720	
C	MPART1	MP000810	
C	MPASIN	MP000860	
C	MPATAN	MP000890	
C	MPBERN	MP000960	
C	MPBESJ	MP010760	
C	MPBES2	MP011900	
C	MPCDM	MP012580	
C	MPCHK	MP013250	
C	MPCIM	MP013760	
C	MPCLR	MP014030	
C	MPCMD	MP014160	
C	MPCMDE	MP014560	
C	MPCMEF	MP014790	
C	MPCMF	MP015500	
C	MPCMI	MP015820	
C	MPCMIM	MP016250	
C	MPCMPA	MP016480	
C	MPCMPI	MP016640	
C	MPCMPR	MP016800	

C	MPCMR	MP016960	MP004660
C	MPCMRE	MP017310	MP004670
C	MPCOMP	MP017530	MP004680
C	MPCOS	MP017880	MP004690
C	MPCOSH	MP018170	MP004700
C	MPCQM	MP018430	MP004710
C	MPCRM	MP018630	MP004720
C	MPDAW	MP019260	MP004730
C	MPDIV	MP019900	MP004740
C	MPDIVI	MP020380	MP004750
C	MPDUMP	MP021510	MP004760
C	MPEI	MP021770	MP004770
C	MPEPS	MP022950	MP004780
C	MPERF	MP023240	MP004790
C	MPERFC	MP023810	MP004800
C	MPERF2	MP024310	MP004810
C	MPERF3	MP024840	MP004820
C	MPERR	MP025390	MP004830
C	MPEUL	MP025580	MP004840
C	MPEXP	MP026340	MP004850
C	MPEXP1	MP027370	MP004860
C	MPEXT	MP028000	MP004870
C	MPGAM	MP028240	MP004880
C	MPGAMQ	MP029050	MP004890
C	MPGCD	MP030300	MP004900
C	MPHANK	MP030560	MP004910
C	MPIN	MP031390	MP004920
C	MPINE	MP032570	MP004930
C	MPLI	MP033020	MP004940
C	MPLN	MP033380	MP004950
C	MPLNGM	MP033940	MP004960
C	MPLNGS	MP034810	MP004970
C	MPLNI	MP035580	MP004980
C	MPLNS	MP036630	MP004990
C	MPL235	MP037300	MP005000
C	MPMAX	MP037830	MP005010
C	MPMAXR	MP037950	MP005020
C	MPMIN	MP038110	MP005030
C	MPMINR	MP038230	MP005040
C	MPMLP	MP038390	MP005050
C	MPMUL	MP038490	MP005060
C	MPMULI	MP039240	MP005070
C	MPMULQ	MP039340	MP005080
C	MPMUL2	MP039600	MP005090
C	MPNEG	MP040480	MP005100
C	MPNZR	MP040560	MP005110
C	MPOUT	MP041370	MP005120
C	MPOUTE	MP041520	MP005130
C	MPOUT2	MP041840	MP005140
C	MPOVFL	MP043230	MP005150
C	MPPACK	MP043440	MP005160
C	MPPI	MP043710	MP005170
C	MPPIGL	MP043970	MP005180
C	MPPOLY	MP044420	MP005190
C	MPPWR	MP044700	MP005200
C	MPPWR2	MP045130	MP005210
C	MPQPWR	MP045440	MP005220
C	MPREC	MP046170	MP005230
C	MPROOT	MP046980	MP005240
C	MPSET	MP048040	MP005250
C	MPSIN	MP048810	MP005260
C	MPSINH	MP049480	MP005270
C	MPSIN1	MP049860	MP005280
C	MPSQRT	MP050440	MP005290
C	MPSTR	MP050680	MP005300
C	MPSUB	MP050890	MP005310
C	MPTAN	MP050980	MP005320
C	MPTANH	MP051610	MP005330
C	MPUNFL	MP052000	MP005340
C	MPUNPK	MP052150	MP005350
C	MPZETA	MP052440	MP005360
C	MP40D	MP053590	MP005370
C	MP40E	MP053760	MP005380
C	MP40F	MP053900	MP005390
C	MP40G	MP054080	MP005400
C	TEST	MP054220	MP005410

```

C TESTV MP056030 MP005420
C TEST2 MP057940 MP005430
C TIMEMP MP064140 MP005440
C MP005450
C $$ ***** EXAMPLE ***** MP005460
C MP005470
C THIS PROGRAM COMPUTES PI AND EXP(PI*SQRT(163/9)) TO 100 MP005480
C DECIMAL PLACES, AND EXP(PI*SQRT(163)) TO 90 DECIMAL PLACES, MP005490
C AND WRITES THEM ON LOGICAL UNIT 6. EXECUTION MP005500
C TIME ON A UNIVAC 1108 (WITH FORTRAN SEID) IS 1.051 SECONDS. MP005510
C MP005520
C TO RUN EXAMPLE THE FOLLOWING MP ROUTINES ARE REQUIRED - MPABS, MP005530
C MPADD, MPADDI, MPADD2, MPADD3, MPART1, MPCHK, MPCIM, MPCLR, MPCMF, MP005540
C MPCMI, MPCMPR, MPCMR, MPCOMP, MPCQM, MPCRM, MPDIVI, MPERR, MP005550
C MPEXP, MPEXP1, MPGCD, MPLNL, MPL235, MPMAXR, MPMLP, MPMUL, MP005560
C MPMULI, MPMULQ, MPMUL2, MPNZR, MPOUT, MPOUT2, MPOVFL, MPPI, MP005570
C MPPWR, MPQPWR, MPREC, MPROOT, MPSET, MPSTR, MPSTR, MPUNFL. MP005580
C MP005590
C CORRECT OUTPUT (EXCLUDING HEADINGS) IS AS FOLLOWS MP005600
C MP005610
C 3.14159265358979323846264338327950288419716939937510 MP005620
C 58209749445923078164062862089986280348253421170680 MP005630
C 640320.00000000060486373504901603947174181881853947577148 MP005640
C 57603665918194652218258286942536340815822646477590 MP005650
C 262537412640768743.999999999925007259719818568887935385633733699086 MP005660
C 2707537410378210647910118607312951181346 MP005670
C MP005680
C CERTAIN PARAMETERS AND WORKING SPACE IN COMMON. MP005690
C COMMON B, T, M, LUN, MXR, R MP005700
C MP005710
C MPEXP REQUIRES 4T+10 WORDS AND WE HAVE T .LE. 62 IF WORDLENGTH MP005720
C AT LEAST 16 BITS, SO 4T+10 .LE. 258. DIMENSIONS CAN BE REDUCED MP005730
C IF WORDLENGTH IS GREATER THAN 16 BITS. MP005740
C INTEGER B, T, R(258) MP005750
C MP005760
C VARIABLES NEED T+2 .LE. 64 WORDS AND ALLOW 110 CHARACTERS FOR MP005770
C DECIMAL OUTPUT MP005780
C INTEGER PI(64), X(64), C(110) MP005790
C MP005800
C CALL MPSET TO SET OUTPUT LOGICAL UNIT = 6 AND EQUIVALENT MP005810
C NUMBER OF DECIMAL PLACES TO AT LEAST 110. THE LAST TWO MP005820
C PARAMETERS ARE THE DIMENSIONS OF PI (OR X) AND R. MP005830
C CALL MPSET (6, 110, 64, 258) MP005840
C MP005850
C COMPUTE MULTIPLE-PRECISION PI MP005860
C CALL MPPI(PI) MP005870
C MP005880
C CONVERT TO PRINTABLE FORMAT (F110.100) AND WRITE MP005890
C CALL MPOUT (PI, C, 110, 100) MP005900
C WRITE (LUN, 10) B, T, C MP005910
C 10 FORMAT (32H1EXAMPLE OF MP PACKAGE, BASE =, I9, MP005920
C $ 12H, DIGITS =, I4 /// 11H PI TO 100D // MP005930
C $ 11X, 60A1 / 21X, 50A1) MP005940
C MP005950
C SET X = SQRT(163/9), THEN MULTIPLY BY PI MP005960
C CALL MPQPWR (163, 9, 1, 2, X) MP005970
C CALL MPMUL (X, PI, X) MP005980
C MP005990
C SET X = EXP(X) MP006000
C CALL MPEXP (X, X) MP006010
C MP006020
C CONVERT TO PRINTABLE FORMAT AND WRITE MP006030
C CALL MPOUT (X, C, 110, 100) MP006040
C WRITE (LUN, 20) C MP006050
C 20 FORMAT (/ 28H EXP(PI*SQRT(163/9)) TO 100D // MP006060
C $ 11X, 60A1 / 21X, 50A1) MP006070
C MP006080
C SET X = X**3 = EXP(PI*SQRT(163)) MP006090
C CALL MPPWR (X, 3, X) MP006100
C MP006110
C WRITE IN FORMAT F110.90 MP006120
C CALL MPOUT (X, C, 110, 90) MP006130
C WRITE (LUN, 30) C MP006140
C 30 FORMAT (/ 25H EXP(PI*SQRT(163)) TO 90D // MP006150
C $ 1X, 70A1 / 21X, 40A1) MP006160
C STOP MP006170
C END MP006180

```

REMARK ON ALGORITHM 524

MP, A Fortran Multiple-Precision Arithmetic Package [A1]
 [R.P. Brent, *ACM Trans. Math. Software* 4, 1 (March 1978), 71-81]

R.P. Brent [Recd 7 Aug. 1978 and 6 Dec. 1978]
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A new version of the Fortran multiple-precision arithmetic package MP, which is described in [1] and given as ACM Algorithm 524, is now available from the ACM Algorithms Distribution Service. The new version may be used with the Augment preprocessor [4], and the necessary interface routines and description deck, described in [2], are supplied. The MP Users' Guide (also supplied with the package) has been revised to describe the Augment interface routines and the use of MP via Augment.

The new version incorporates faster algorithms for the exponential integral and Euler's constant [3]. Consequently, the TEST2 program described in [1] now runs about 10 percent faster.

In versions of MP dated June 7, 1978, and earlier (including Algorithm 524), there may be an error in the least significant digit when multiple-precision numbers with the same sign and exponent are added using subroutine MPADD. To correct this, change .LT. to .LE. in line MP007390 of MPADD3.

ACKNOWLEDGMENT

I am indebted to John P. Jeter of the University of South West Louisiana for finding both the error in MPADD3 and the required correction.

REFERENCES

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2. BRENT, R.P., HOOPER, J.A., AND YOHE, J.M. An Augment interface for Brent's multiple-precision arithmetic package. Mathematics Res. Ctr., U. of Wisconsin, Madison, Wis., Aug. 1978.
3. BRENT, R.P., AND MCMILLAN, E.M. Some new algorithms for high-precision computation of Euler's constant. To appear in *Mathematics of Computation*.
4. CRARY, F.D. A versatile precompiler for nonstandard arithmetics. *ACM Trans. Math. Software* 5, 2 (June 1979), 204-217.

ALGORITHM

[Only summary information of the revised version of the algorithm is printed here. This listing is available from the ACM Algorithms Distribution Service (see inside back cover for order form) and will be supplied to those requesting Algorithm 524.]

NAME(n): indicates a Fortran module from the MP package with n records
 NAME^B(n): indicates "NAME" is included as an example
 NAME^T(n): indicates "NAME" is part of the test package
 NAME^U(n): indicates a listing of the user guide
 NAME^A(n): indicates an Augment description deck and Jacobi program using it

Contents: EXAMPLE^T(111), MPABS(8), MPADD(8), MPADDI(16), MPADDQ(12), MPADD2(60), MPADD3(90), MPART1(48), MPASIN(38), MPATAN(62), MPBASA(9), MPBASB(13), MPBERN(115), MPBESJ(114), MPBES2(68), MPCAM(44), MPCDM(67), MPCHK(51), MPCIM(26), MPCLR(13), MPCMD(40), MPCMDE(23), MPCMEF(71), MPCMF(32), MPCMI(43), MPCMIM(23), MPCMPA(16), MPCMPI(16), MPCMPR(16), MPCMR(35), MPCMRE(22), MPCOMP(35), MPCOS(29), MPCOSH(26), MPCQM(20), MPCRM(63),

MPDAW(64), MPDGA(10), MPDGB(21), MPDIGA(9),
MPDIGB(14), MPDIV(48), MPDIVI(113), MPDUMP(26),
MPEI(156), MPEPS(29), MPEQ(7), MPERF(57), MPERFC(50),
MPERF2(53), MPERF3(55), MPERR(19), MPEUL(76),
MPEXP(103), MPEXPA(11), MPEXPB(24), MPEXP1(63),
MPEXT(24), MPGAM(81), MPGAMQ(125), MPGCD(26),
MPGCD A(78), MPGCD B(38), MPGE(7), MPGT(7),
MPHANK(83), MPIN(118), MPINE(45), MPINF(27),
MPINIT(39), MPIO(19), MPKSTR(15), MPLE(7), MPLI(36),
MPLN(56), MPLNGM(87), MPLNGS(77), MPLNI(105),
MPLNS(67), MPLT(7), MPL235(53), MPMAX(12),
MPMAXR(16), MPMEXA(9), MPMEXB(18), MPMIN(12),
MPMINR(16), MPMLP(10), MPMUL(75), MPMULI(10),
MPMULQ(26), MPMUL2(88), MPNE(7), MPNEG(8),
MPNZR(81), MPOUT(15), MPOUTE(32), MPOUTF(25),
MPOUT2(139), MPOVFL(21), MPPACK(27), MPPI(26),
MPPIGL(45), MPPOLY(28), MPPWR(43), MPPWR2(31),
MPQPWR(73), MPREC(81), MPROOT(102), MPSET(77),
MPSIGA(7), MPSIGB(24), MPSIN(67), MPSINH(38),
MPSIN1(58), MPSQRT(24), MPSTR(21), MPSUB(9),
MPTAN(63), MPTANH(39), MPUNFL(15), MPUNPK(29),
MPUPK(45), MPZETA(115), MP40D^T(17), MP40E^T(14),
MP40F^T(18), MP40G^T(15), TEST^T(181), TESTV^T(191),
TEST2^T(620), TIMEMP^T(28), GUIDE^U(1719), AUGDECK^A(159)

ALGORITHM 525

ADAPT, Adaptive Smooth Curve Fitting [E2]

JOHN R. RICE
Purdue University

Key Words and Phrases: piecewise polynomial approximation, adaptive curve fitting, adaptive approximation, Hermite interpolation, function approximation
CR Categories: 5.13
Language: Fortran

DESCRIPTION

1. Introduction and Background

The basic objectives of this function approximation algorithm are: speed, reliability, and smoothness. Algorithms already exist with any two of these three properties. Speed requires that the work be proportional to the length of the curve (for fixed accuracy and more or less uniform complexity of the function). Reliability requires that curves with singularities or near singularities, oscillations, and other complex behavior be handled. Smoothness (number of continuous derivatives) of the approximation obtained is input to the algorithm. The author believes that adaptive piecewise polynomial algorithms offer the best hope for such algorithms with nonadaptive piecewise polynomial schemes or rational function approximation as the only serious competitors. Nonadaptive schemes do not cope efficiently with functions having very nonuniform behavior (singularities in slopes, for example) and the work for rational approximation probably increases faster than linearly with the length of the curve.

The theoretical background for this algorithm is provided by Rice [1] and the references cited there. This may be summarized by saying that for given fixed degree n of the pieces the error of the best piecewise polynomial approximation behaves like k^{-n} where k is the number of pieces. This theoretical result applies to a broad class of functions which includes everything that conceivably could arise in applications. A class of adaptive algorithms has been found which chooses the knots so that the error behaves like k^{-n} even though the best approximation is not obtained. These algorithms apply to any piecewise smooth function with a finite number of "algebraic" singularities, i.e. to any function $F(x)$ which behaves like $a + b(x - s)^\alpha$ near the singularity s . The exponent α must give a finite value for the norm of $F(x)$ specified for ADAPT to use, i.e. $\alpha > -\frac{1}{2}$ for least squares, $\alpha > 0$ for uniform approximation.

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The key ingredients in these algorithms are a local approximation operator, a local error estimator, and a data structure for the subintervals generated. ADAPT uses Hermite interpolation at subinterval endpoints plus ordinary interpolation in between if needed. Error estimates are made by a simple Gauss quadrature formula and a stack is used for the intervals. This algorithm is probably the simplest choice from those currently known to give the optimal convergence rate. A more detailed description of ADAPT is given in Section 5. Note that the nature of ADAPT *requires* that $F(x)$ and its derivatives be available for arbitrary x (see [2] for guidance on numerically estimating derivatives) and thus ADAPT is not directly applicable to discrete data sets.

Section 2 presents some remarks about the algorithm components; most of them incorporate rather standard methods. Section 3 discusses the use of ADAPT (input/output, role of the arguments CHARF = characteristic length of $F(x)$ and EDIST = error distribution type, and portability). Section 4 describes the extensive testing performed and summarizes the algorithm properties observed in [2]. It also briefly describes the driver program and 20 functions with derivatives used for testing ADAPT which are applicable to similar algorithms. The final section has a very high level description of the algorithm; ADAPT itself is available through the ACM Algorithms Distribution Service.

2. The Principal Algorithm Components

The algorithm ADAPT is to approximate the function $F(x)$ on the interval $[A,B]$ to within an accuracy ACCUR by piecewise polynomials of degree DEGREE with a number SMOOTH of continuous derivatives. The error is measured in the L_p -norm with p set by NORM. Other input is the characteristic length CHARF of F , printed output level LEVEL, the error distribution type EDIST, and the number NBREAK of breakpoints (plus information about the breakpoints if $NBREAK > 0$). Breakpoints are points where there are breaks in the curve or its derivatives: knots are placed there and special calculations made. The output is the array XKNOTS of knots, the coefficients COEFS (relative to the knot locations) of the polynomial pieces, the estimated error ERROR, and number KNOTS of knots. The approximation obtained is automatically available as the FUNCTION subprogram PPOLY.

2.1 Data Structure and Discard Procedure for Subintervals. ADAPT generates a set of subintervals of $[A,B]$ which are maintained in a stack with the leftmost on the top. Since subintervals are created by halving, the maximum size of the stack is limited by the machine word length.

Intervals are discarded whenever the estimated error on an interval is small enough. This decision made in the subprogram CHECK is somewhat subtle and three strategies are provided. This question is discussed later in some detail in connection with the argument EDIST which indicates the strategy selected by the user. Two subprograms, PUT and TAKE, are used for access to the interval stack.

2.2 Hermite Interpolation. The values of DEGREE and SMOOTH are specified by the user with $DEGREE \geq 2 * SMOOTH + 1$. The polynomial pieces are determined by interpolating F and SMOOTH derivatives of F at each endpoint of a subinterval plus the value of F at $DEGREE - 2 * SMOOTH - 1$ other points. During the computation the polynomial pieces are represented by divided differences computed in NEWTON and used by POLYDD. Once a polynomial piece is accepted for the final approximation PTRANS transforms its representation into powers with origin at the left endpoint of the subinterval. The transformation is accomplished by repeated synthetic division. The subprogram COMPUT controls the computation of a polynomial piece as well as the error estimation.

2.3 The Breakpoint Mechanism. It is sometimes very useful for a user to be able to specify breaks in some derivative at certain points. Most commonly one has a known or desired jump in the slope at a given point and ADAPT allows this via NBREAK and associated arguments XBREAK, DBREAK, BLEFT, BRIGHT which specify the exact nature of the breakpoint. This rather straightforward facility is implemented primarily in TAKE with some impact on COMPUT.

2.4 *Error Estimation.* The L_P -norm of the local error is estimated in ERRINT by a 4-point Gauss quadrature for $(F - \text{POLYDD})^{**}P$ on the subinterval under consideration. Special code is used for $P = \text{infinity}$ (minimax approximation). The global error estimate is built up in PUT by appropriately combining the local error estimates.

2.5 *Fatal Errors.* The algorithm normally terminates when the stack is empty. The stack should not overflow but might do so at very strong singularities which cause the algorithm to want to operate at accuracies inconsistent with the machine word length. This overflow has not occurred in the testing so far, but if it does a message such as the following is printed:

```
INTERVAL DIVIDED TOO MUCH, EXCEEDED LIMIT 50 ON INTERVAL STACK
AT
17923.12345678    17923.12345687
INTERVAL DISCARDED AND COMPUTATION CONTINUED
```

This message may be suppressed by setting $\text{LEVEL} = -1$, but note that the computation is allowed to proceed on the conjecture that this situation is not truly fatal.

A strong singularity has been observed to cause another situation indicated by a message like

```
GOT SHORT INTERVAL **** 3210.12345678  3210.12345679 **** DISCARD IT
```

which may also be suppressed by setting $\text{LEVEL} = -1$. Experiments indicate that the algorithm will recover and produce satisfactory results provided it can discard (and ignore) enough short intervals to get out of the region where the machine word length is inadequate. It often cannot get out of this region before exceeding run time limits or, more likely, the limit on the number of knots. The variable BUFFER in PUT governs short interval detection.

The arrays XKNOTS and COEFS are passed to ADAPT with variable dimensions KDIMEN and NDIMEN. If the number of knots computed exceeds KDIMEN then a fatal error message is printed and the computation aborted. This message cannot be suppressed. Checks are made on various input parameters by SETUP and inconsistent or impossible input leads to fatal error messages and a RETURN without any computation.

3. Algorithm Usage

3.1 *User Interface.* The input to ADAPT is via the COMMON block INPUTZ except for the function F. This is the most convenient for extensive use of ADAPT but not for occasional use. A subroutine PPFIT4 is provided which has all input as arguments. Entry points PPFIT1, PPFIT2, and PPFIT3 in PPFIT4 have fewer (10, 12, 15, respectively) arguments than PPFIT4 (21 arguments). Due to the variable nature of entry point implementations, these are only indicated by COMMENT cards and local modifications are needed to activate these features.

The basic output is the arrays XKNOTS and COEFS which are arguments to ADAPT (and also the PPFIT subroutines) and the numbers KNOTS and ERROR in the COMMON block RESULTZ. The PPFIT routines have KNOTS and ERROR as arguments. In addition the FUNCTION subprogram PPOLY (T, XKNOTS, COEFS, KDIMEN, NDIMEN) returns the value at the point T of the most recently computed approximation. It is automatically available to the user.

There are six levels (-1, 0, 1, 2, 3, 4) of printed output available. $\text{LEVEL} = -1$ only provides fatal error messages; $\text{LEVEL} = 0$ also provides "semifatal" error messages plus 1 line of output; $\text{LEVEL} = 1$ provides a printout of the input and the approximation obtained; $\text{LEVEL} = 2$ provides a condensed trace of the computation and the last two levels are only useful for debugging program modifications.

3.2 *The Characteristic Length of F, CHARF.* The correct operation of ADAPT depends on certain estimates being accurate which, in turn, depend on the relevant subintervals being small enough. Specifically, the sampling that ADAPT does of

F must reveal the nature of F and not allow any significant features to go undetected. The nature of ADAPT is such that for reasonable functions it is very unlikely to miss significant features without using CHARF at all. However, knowing how ADAPT works, one can readily construct examples where it will fail unless CHARF is set properly. The value of CHARF is an *upper* limit on the size of the subintervals for polynomial pieces and it is to be set so that the 4-point Gauss quadrature formulas are reasonably (but not highly) accurate. This means that if F has some complex behavior on a very short segment, then setting CHARF to, say, half the length of this segment will force ADAPT to detect this behavior.

The argument CHARF is essential to proving ADAPT correct (which has not been attempted) but its practical value is debatable. If F is more or less uniformly complicated then short intervals are needed everywhere. It is extremely unlikely, but not impossible, that ADAPT would be fooled in such a case. If F is very smooth except on a very short segment, then ADAPT may well be fooled and setting CHARF small will avoid this. However, it will also force very small intervals where F is smooth and where they are not needed. Thus high reliability is obtained here at the cost of great inefficiency.

The nature of the approximations computed by ADAPT are such that there is a simple and efficient alternative to setting CHARF small. Let [C,D] be the subinterval of [A,B] where F is rough and let F be smooth elsewhere. One can then approximate F on [A,C], [C,D], and [D,B] independently, and the approximations fit together smoothly at C and D to give a single smooth approximation for the entire interval [A,B].

3.3 *The Error Distribution Type, EDIST.* The least squares error for the approximation S(x) to F(x) is

$$E(A,B) = \left[\int_A^B (F(x) - S(x))^2 dx \right]^{1/2}.$$

To make $E(A,B) < .01$ is equivalent to making $E(A,B)^2 < .0001$, and for any L_p -norm the program actually operates with $ACCUR^p$. Suppose $A=0, B=1$, then we can achieve $E(0,1)^2 < .0001$ by achieving $E(0,1/2)^2 < .0001/2$ and $E(1/2, 1)^2 < .0001/2$ since $E(A,B)^2$ is simply additive over intervals. This approach is called *proportional error distribution*, i.e. the total error requirement is distributed over the subintervals of [A,B] in proportion to the lengths of the subintervals. This choice is selected by $EDIST=0$ and automatically results in the total error $E(A,B)^2$ less than the specified error $ACCUR^2$.

An alternative approach is to make the errors approximately equal on each of the subintervals independent of their lengths. This is called *fixed error distribution* and is selected by $EDIST=2$. The argument $ACCUR$ is used for each subinterval and if k subintervals are finally used we see that the total error is then

$$[k * ACCUR^p]^{1/p} = \sqrt[p]{k} ACCUR$$

and thus $ACCUR$ is not the specified total error when $EDIST=2$. This alternative is awkward to use because the final approximation error depends on the number of subintervals required which, of course, is unknown until after the approximation is computed. However, for rough or singular F(x) this disadvantage is more than compensated by the superior performance of this error distribution type. This is seen in the theory and verified in actual use.

A compromise approach called *approximate fixed error distribution* is selected by $EDIST=1$. Basically the algorithm keeps a running estimate of the final number of intervals it will use and adjusts the error requirement for subintervals accordingly. This approach is obviously not foolproof, but the testing reported below shows it to be 98 to 99 percent reliable. In any case, the total error actually obtained is available for the user to see and to test. For smooth, uniformly varying F(x) the proportional error distribution gives perfectly satisfactory efficiency.

We observe that for $p = \text{infinity}$ ($NORM=3$) there is no difference between the fixed and proportional error distributions.

3.4 *Portability Considerations.* Considerable pains have been taken to make the Fortran program portable. It is written in a subset of ANSI Fortran specified by PFORT (see [5]) except that four characters are packed per word rather than one as specified by PFORT. The current version is in single precision and specifically tailored to a machine with a long word length (CDC 6000-7000 series computers). Specific directions are given in the comments for changing the precision or to use it on machines with a shorter word length. In particular, all REAL variables are explicitly declared to facilitate the change to DOUBLE PRECISION. The program has been generated by an experimental Fortran converter which automatically produces versions tailored for different machines and precisions. Several of these versions have been produced and run successfully.

4. Algorithm Testing and Verification

This algorithm is based on a method with theoretically known properties. Care was taken to adhere to the requirements of that theory and considerable analysis has been made of various features and parts of the algorithm. However, no formal proof has been attempted since, as with many numerical algorithms, one cannot say a priori what is to be computed. A somewhat related algorithm has been proved correct with a few idealizing assumptions (infinite precision arithmetic and infinite memory, for example); see [3, 4].

Very extensive testing of the algorithm has been performed to see if the theoretical expectations are, in fact, realized by this algorithm. There are some approximation methods where this has not been the case. These tests are discussed at some length in [2] and we summarize them by saying that this algorithm performs as expected from the theory. The results in [2] give many insights into the practical use of this algorithm. About 2500 to 3000 different approximations have been computed in these tests and all have been examined for signs of incorrect performance. A few hundred of these runs were specifically designed to test the validity and correctness of the program.

The code with the ADAPT includes a set of 20 test functions (10 of them param-

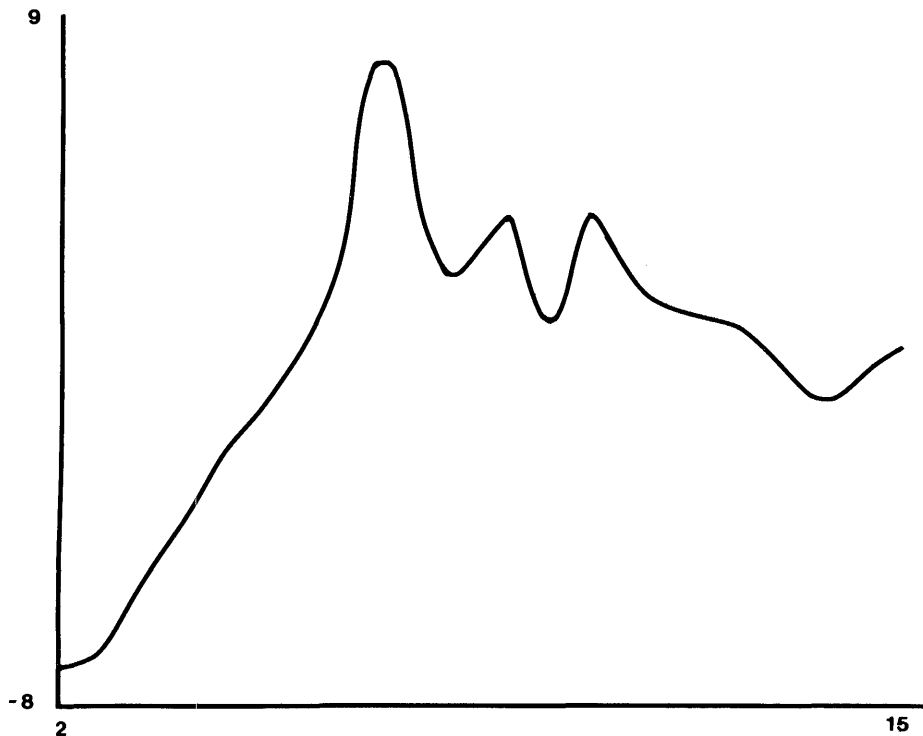


Fig. 1. The complicated function with seven pieces approximated in the second example

terized in various ways) and a driver to exercise ADAPT and measure its performance in various ways. These are included because it is felt they will be useful for testing other function approximation programs. We note that the 20 test functions are chosen to present various features of the approximation problem and we found it amazingly difficult to get the first few derivatives of these functions computed correctly.

We give sample output for two functions, one easy and one difficult. The first is $F(x) = 1./(1+(x-2.5)^4)$ with LEVEL=2 output given. The second function is shown in Figure 1 with LEVEL = 0 output given. In the first example the first and last three lines of output are from the test driver and the specific information about the problem is printed by ADAPT. For the second example the approximation is on [2,15] for polynomial degree 6 with 1 continuous derivative; the requested and estimated accuracies are .001 and .00057 in the minimax norm; proportional error distribution was used with CHARF=4. This example is function 18 of the set of test functions and one of its parameters, P8F, has been set to 1.6 here instead of 0.6 as given in the DATA for F(x). The output has been reformatted to accommodate the narrowed page width here.

```
*****ADAPT FOR FUNCTION 4 = HUMP AT 2.5 RECIPROCAL OF QUARTIC
PIECEWISE POLYNOMIAL APPROXIMATION ON INTERVAL 1.0000E+00 5.0000E+00
OF DEGREE 5 WITH 2 CONTINUOUS DERIVATIVES
ACCURACY REQUESTED IS 1.0000E-03 MEASURED BY LEAST SQUARES
OTHER INPUT/DEFAULT VARIABLES ARE FOSCIL = 4.0000E+00
EMEAS = 2.0000E+00 ---- PROPORTIONAL ERROR DISTRIBUTION
KNOT 2 AT 1.50000E+00, LOCAL-GLOBAL ERRORS = 2.6008E-09 5.0998E-05
KNOT 3 AT 2.00000E+00, LOCAL-GLOBAL ERRORS = 3.5626E-08 1.9552E-04
KNOT 4 AT 2.50000E+00, LOCAL-GLOBAL ERRORS = 2.5115E-08 2.5168E-04
KNOT 5 AT 3.00000E+00, LOCAL-GLOBAL ERRORS = 2.5115E-08 2.9742E-04
KNOT 6 AT 3.50000E+00, LOCAL-GLOBAL ERRORS = 3.5626E-08 3.5225E-04
KNOT 7 AT 4.00000E+00, LOCAL-GLOBAL ERRORS = 2.6008E-09 3.5593E-04
KNOT 8 AT 5.00000E+00, LOCAL-GLOBAL ERRORS = 2.4834E-09 3.5940E-04
```

```
--- ADAPTIVE PIECEWISE POLYNOMIAL APPROXIMATION OF DEGREE 5 WITH
2 CONTINUOUS DERIVATIVES NEEDED 8 KNOTS FOR ERROR = 3.5940E-04
KNOT LOCATION      X-POWER COEFFICIENTS RELATIVE TO KNOT LOCA-
                   TIONS
1  1.000000000000E+00  1.649484536082E-01
                   3.673078966947E-01
                   4.506148423367E-01
                   2.850451366482E-01
                   5.027123624777E-01
                   -9.058008728205E-01
2  1.500000000000E+00  5.000000000000E-01
                   1.000000000000E+00
                   5.000000000000E-01
                   -1.032973743131E+00
                   -2.117239975575E+00
                   2.484021982495E+00
3  2.000000000000E+00  9.411764705882E-01
                   4.429065743945E-01
                   -1.120293099939E+00
                   7.978831671079E-01
                   7.848565031541E-01
                   -1.003053124363E+00
4  2.500000000000E+00  1.000000000000E+00
                   0
                   0
                   1.400366374921E-01
                   -1.722776307754E+00
                   1.003053124363E+00
5  3.000000000000E+00  9.411764705882E-01
                   -4.429065743945E-01
                   -1.120293099939E+00
                   -9.426012619580E-01
                   4.092814980663E+00
```

```

-2.484021982495E+00
6 3.500000000000E+00 5.000000000000E-01
-1.000000000000E+00
5.000000000000E-01
9.740323204476E-01
-1.761789819574E+00
9.058008728205E-01
7 4.000000000000E+00 1.649484536082E-01
-3.673078966947E-01
4.506148423367E-01
-3.547233211608E-01
1.658371414434E-01
-3.440822109296E-02

```

ADAPT USED 117 FUNCTION VALUES FOR ERRORS
 SPECIFIED = 1.00000E-03 ESTIMATED BY ADAPT = 3.59399E-04
 INDEPENDENT CHECK = 3.05041E-04 TIME USED = .17200 SECONDS

*****ADAPT FOR FUNCTION 18 = COMPLICATED FUNCTION WITH 7 PIECES
 -----ADAPTIVE PIECEWISE POLYNOMIAL APPROXIMATION OF DEGREE 6 WITH
 1 CONTINUOUS DERIVATIVES NEEDED 30 KNOTS FOR ERROR = 5.7301E-04

5. High Level Expression of the Algorithm ADAPT

The following description of ADAPT indicates its basic structure and methods of computation:

```

PROCEDURE PPFIT - WITH ARGUMENTS =
  F      .FUNCTION TO FIT      LEVEL      .OUTPUT LEVEL -1 TO 2
  A,B    .INTERVAL ENDPTS      CHARF      .LIMIT ON PIECES LENGTH
  ACCUR  .ACCURACY DESIRED     EDIST    .TYPE OF ERROR CONTROL
  DEGREE .POLYNOMIAL DEGREE   NBREAK   .NUMBER OF SPECIFIED
  SMOOTH .NO. CONT. DERIVS    .NUMBER OF BREAK POINTS IN FIT.
  NORM   .MEAS. OF L-P ERROR  HAS RELATED ARGUMENTS
        UP IN (0, INFINITY)   GIVING DETAILS.

  OUTPUT = KNOTS.              .NUMBER OF KNOTS
  ERROR.                       .ERROR ACHIEVED
  XKNOTS(K), K=1 TO KNOTS     .KNOT LOCATIONS
  COEFS(K,N), K=1 TO KNOTS   .POWER COEF. RELATIVE
        N=1 TO DEGREE+1     .TO THE KNOT LOCATIONS

```

*** THIS PROGRAM MERELY COMPUTES A FEW DEFAULT VALUES
 AND PUTS VARIABLES IN COMMON BLOCKS, ETC

```

  CALL ADAPT - TO DO THE APPROXIMATION
END PPFIT
SUBPROGRAM ADAPT
  CALL SETUP - CHECK INPUT, INITIALIZE THINGS, PRINT PROBLEM

  *** LOOP OVER PROCESSING INTERVALS ***

  CALL TAKE -- AN INTERVAL OFF THE STACK
  CALL COMPUT -- AN APPROX ON THIS INTERVAL
  CALL CHECK -- FOR DISCARDING OR DIVIDING INTERVAL
  CALL PUT -- NEW INTERVALS ON STACK, UPDATE ALGORITHM STATUS
  CALL TERMIN -- TEST FOR FINISH, PRINT INTERMEDIATE OUTPUT
  IF NOT FINISHED - REPEAT LOOP

  CALL SUMMARY -- FOR FINAL OUTPUT
END ADAPT

SUBPROGRAM SETUP
  SET LIMITS ON COMPUTATION PARAMETERS
  CHECK ALL INPUT DATA
  INITIALIZE VARIABLES AND INTERVAL STACK
  PRINT PROBLEM STATEMENT
END SETUP

SUBPROGRAM TAKE
  CHECK FOR BREAK POINT IN TOP INTERVAL

```

```

    IF SO - ADJUST XKNOTS TO MAKE IT A PARTITION POINT
    ELSE - DO NOTHING
END TAKE

SUBPROGRAM PUT
CHECK FOR DISCARDING INTERVAL
  IF SO - UPDATE ERROR ESTIMATE
    ADJUST STACK
    CALL PTRANS - TO OBTAIN COEFS FOR THIS INTERVAL
    UPDATE XKNOTS AND COEFS
  ELSE - SUBDIVIDE INTERVAL AND PLACE 2 NEW ONES ON STACK
    CHECK FOR EXCEEDING MAX STACK SIZE OR OBTAINING
    AN INTERVAL WHICH IS TOO SHORT. SUCH SHORT INTERVALS
    ARE DISCARDED WITHOUT REGARD TO ERROR CONTROL POLICY
    AND WITH MESSAGE
END PUT

SUBPROGRAM PTRANS - OF PUT
CHANGES POLYNOMIAL REPRESENTATION FROM NEWTON DIVIDED
DIFFERENCE FORM TO POWER FORM WITH ORIGIN SHIFTED TO THE
XKNOT VALUE ON LEFT OF INTERVAL. USES SYNTHETIC DIVISION
END PTRANS

SUBPROGRAM COMPUT
OBTAIN - VALUES OF F AND DERIVATIVES. MAKE ADJUSTMENTS
        IF A BREAK POINT IS INVOLVED
CALL NEWTON - FOR DIVIDED DIFFERENCES OF INTERPOLATING
        POLYNOMIAL FOR THIS INTERVAL
CALL ERRINT - TO ESTIMATE LOCAL ERROR + (F(X)-POLYDD(X))**P
        FOR L-P NORM, 0 < P <= INFINITY
END COMPUT

SUBPROGRAM NEWTON - OF COMPUT
BUILD UP TRUE DIVIDED DIFFERENCE TABLE WITH MULTIPLE POINTS
AT THE INTERVAL ENDS PLUS INTERPOLATION POINTS
END NEWTON

SUBFUNCTION POLYDD - OF COMPUT
EVALUATES POLYNOMIAL FROM DIFFERENCE TABLE OUT OF NEWTON
FOR L-P NORM WITH P IN (0, INFINITY)
END POLYDD

SUBPROGRAM ERRINT - OF COMPUT
USES 4-POINT GAUSS QUADRATURE TO ESTIMATE ERROR NORM ON
INTERVAL. SPECIAL COMPUTATION FOR MAX-NORM, P=INFINITY.
END ERRINT

SUBPROGRAM CHECK
USES ERROR DISTRIBUTION TYPE AND CHARF TO DECIDE ON DISCARD
END CHECK

SUBPROGRAM TERMIN
PRINT - INTERMEDIATE OUTPUT, IF ANY REQUESTED BY LEVEL = 2
TEST - FOR TERMINATION EMPTY STACK - NORMAL
        EXCEEDED XKNOTS LIMIT - ABNORMAL
END TERMIN

SUBPROGRAM SUMMARY
LEVEL = -1 NOTHING
        = 0 1 LINE (OUTPUT RETURNED IN COMMON, ARGUMENTS)
        = 1 KNOTS AND COEFFICIENTS
        = 2 DITTO
END SUMMARY

FUNCTION PPOLY
THE PIECEWISE POLYNOMIAL COMPUTED BY PREVIOUS CALL ON PPFIT
END PPOLY

```

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3. RICE, J.R. Parallel algorithms for adaptive quadrature II—Metalgorithm correctness. *Acta Inform.* 4 (1975), 273-285.
4. RICE, J.R. Parallel algorithms for adaptive quadrature III—Program correctness. *ACM Trans. Math. Software* 2, 1 (March 1976), 1-30.
5. RYDER, B.G. The PFORT verifier. *Software—Practice and Experience* 4 (1974), 359-377.

ALGORITHM

[The listing printed here is an abridged version of the complete algorithm, which is available upon request from the ACM Algorithms Distribution Service.]

```

SUBROUTINE ADAPT(F, XKNOTS, COEFS, KDIMEN, NDIMEN)          ADA 10
C                                                         ADA 20
C THIS ALGORITHM COMPUTES A PIECEWISE POLYNOMIAL APPROXIMATION ADA 30
C OF SPECIFIED SMOOTHNESS, ACCURACY AND DEGREE. THE INPUT TO THE ADA 40
C COMPUTATION IS                                         ADA 50
C                                                         ADA 60
C F - FUNCTION BEING APPROXIMATED. IT MUST PROVIDE VALUES OF ADA 70
C DERIVATIVES UP TO THE ORDER OF SMOOTHNESS SPECIFIED FOR ADA 80
C THE APPROXIMATION. THE CALLING SEQUENCE IS F(X,FDERV) AND ADA 90
C FDERV CONTAINS THE DERIVATIVES( SEE CONSTRAINT BELOW) ADA 100
C A,B - THE ENDPOINTS OF THE INTERVAL OF APPROXIMATION ADA 110
C ACCUR - THE ACCURACY REQUIRED FOR THE APPROXIMATION ADA 120
C SMOOTH - THE SMOOTHNESS REQUIRED FOR THE APPROXIMATION ADA 130
C           = 0 MEANS CONTINUOUS ADA 140
C           = 1 MEANS CONTINUOUS SLOPE ADA 150
C           = 2 MEANS CONTINUOUS SECOND DERIVATIVE, ETC. ADA 160
C DEGREE - THE DEGREE OF THE POLYNOMIAL PIECES. ADA 170
C           MUST HAVE DEGREE GT 2*SMOOTH ADA 180
C                                                         ADA 190
C ***** SECONDARY INPUT - ITEMS WITH DEFAULT VALUES POSSIBLE ADA 200
C CHARF - CHARACTERISTIC LENGTH OF THE FUNCTION F(X). PIECES ARE NOT ADA 210
C LONGER THAN THIS LENGTH. ADA 220
C           DEFAULT=(B-A) IF DEGREE GT 1, ELSE (B-A)/3 ADA 230
C NORM - NORM TO MEASURE THE APPROXIMATION ERROR ADA 240
C           = 1 L1 APPROXIMATION (LEAST DEVIATIONS) ADA 250
C           = 2 L2 APPROXIMATION (LEAST SQUARES) ADA 260
C           = 3 TCHEBYCHEFF (MINIMAX) APPROXIMATION ADA 270
C           --P (NEGATIVE VALUE) GENERAL LP APPROXIMATION ADA 280
C           DEFAULT= 2 ADA 290
C NBREAK - NUMBER OF SPECIAL BREAK POINTS IN THE APPROXIMATION. ADA 300
C ASSOCIATED INPUT VARIABLES ARE ADA 310
C XBREAK(J) - LOCATION OF BREAK POINTS ADA 320
C DBREAK(J) - DERIVATIVE BROKEN AT XBREAK ADA 330
C BLEFT(J) - VALUE FROM LEFT FOR DBREAK DERIVATIVE ADA 340
C BRIGHT(J) - - - RIGHT - - - ADA 350
C           DEFAULT = 0 ADA 360
C LEVEL - LEVEL OF OUTPUT FROM ADAPT ADA 370
C           = -1 NO OUTPUT AT ALL EXCEPT FOR FATAL INPUT ERRORS ADA 380
C           = 0 ERROR CONDITIONS NOTED, FINAL SUMMARY ADA 390
C           = 1 PRINT THE APPROXIMATION OUT ADA 400
C           = 2 DETAILS OF THE COMPUTATION ADA 410
C           = 3 DEBUG OUTPUT, = 4 LOTS OF DEBUG OUTPUT ADA 420
C           DEFAULT = 0 ADA 430
C EDIST - SWITCH TO CHANGE FROM PROPORTIONAL ERROR DISTRIBUTION ADA 440
C TO FIXED DISTRIBUTION. THIS IS PRIMARILY OF USE IN ADA 450
C APPROXIMATION OF FUNCTIONS WITH SINGULARITIES. ONE SHOULD ADA 460
C USE NORM = 1. OR SO IN SUCH CASES ADA 470
C           = 0 PROPORTIONAL DISTRIBUTION ADA 480
C           = 1 APPROXIMATE FIXED ERROR DISTRIBUTION ADA 490
C           ATTEMPTS TO ACHIEVE SPECIFIED ACCURACY VALUE ACCUR ADA 500
C           = 2 TRUE FIXED ERROR DISTRIBUTION ADA 510
C                                                         ADA 520
C THE OUTPUT OF THE COMPUTATION CONSISTS OF 3 PARTS, EACH RETURNED ADA 530
C TO THE USER IN A DIFFERENT WAY. THEY ARE ADA 540
C                                                         ADA 550
C XKNOTS,COEFS - ARRAYS DEFINING THE PIECEWISE POLYNOMIAL RESULT. ADA 560
C XKNOTS(K) = KNOTS OF THE APPROXIMATION ( K = 1 TO KNOTS) ADA 570
C           THE LAST ONE IS RIGHT END POINT OF INTERVAL ADA 580
C COEFS(K,N) = COEFFICIENT OF (X - XKNOT(K))**(N-1) IN THE ADA 590

```

```

C          INTERVAL XKNOT(K) TO XKNOT(K+1)          ADA 600
C          K = 1 TO KNOTS-1 AND N = 1 TO DEGREE+1  ADA 610
C          THESE ARRAYS ARE PASSED AS ARGUMENTS SO AS TO USE VARIABLE ADA 620
C          DIMENSIONS.                               ADA 630
C          KDIMEN - DIMENSION USED FOR XKNOTS IN CALLING PROGRAM ADA 640
C          NDIMEN - COEFS IS DECLARED COEFS(KDIMEN,NDIMEN) IN THE ADA 650
C          CALLING PROGRAM.                          ADA 660
C          ***** NOTE ***** SEVERAL SMALL ARRAYS HERE HAVE FIXED ADA 670
C          DIMENSIONS THAT LIMIT DEGREE AND THUS NDIMEN ADA 680
C          SHOULD NOT EXCEED THIS LIMIT (CURRENTLY = 12) ADA 690
C          ADA 700
C          PPOLY - THE PIECEWISE POLYNOMIAL APPROXIMATING FUNCTION. ADA 710
C          THIS FUNCTION SUBPROGRAM IS AVAILABLE TO THE USER AT THE ADA 720
C          COMPLETION OF ADAPT.                       ADA 730
C          ADA 740
C          RESULTZ - A LABELED COMMON BLOCK WITH ERROR,KNOTS IN IT ADA 750
C          KNOTS - NUMBER OF KNOTS OF PPOLY           ADA 760
C          ERROR - ACCURACY OF PPOLY AS ESTIMATED BY ADAPT ADA 770
C          ADA 780
C          ***** DIMENSION CONSTRAINTS ***** ADA 790
C          MAXKNT - MAX NUMBER OF KNOTS TAKEN FROM USER VIA KDIMEN ADA 800
C          ARRAYS WITH THIS DIMENSION                 ADA 810
C          COEFS XKNOTS                               ADA 820
C          MAXPAR - MAX NUMBER OF PARAMETERS PER INTERVAL ( = 12 CURRENTLY ) ADA 830
C          USER PROVIDED NDIMEN MUST HAVE NDIMEN LE MAXPAR ADA 840
C          MUST HAVE DEGREE + 1 LE MAXPAR            ADA 850
C          ARRAYS WITH THIS DIMENSION (OR RELATED VALUES ) ADA 860
C          D      DDTMP FDERVL FDERVR FDUMB FACTOR ADA 870
C          FINTRP FLEFT FRIGHT POWERS XTEMP XINTRP XDD ADA 880
C          ***** NOTE ***** MAXPAR ALSO AFFECTS ARGUMENT FDERV ADA 890
C          OF FUNCTION F. FDERVL, FDERVR ARE ALSO INVOLVED. ADA 900
C          SHOULD DECLARE FDERV OF SIZE 6 IN F TO BE SAFE. ADA 910
C          MAXAUX - MAXIMUM NUMBER OF AUXILIARY INPUT( = 20 NOW ). ARRAYS ADA 920
C          XBREAK DBREAK BLEFT BRIGHT              ADA 930
C          MAXSTK - MAX SIZE OF ACTIVE INTERVAL STACK ADA 940
C          MIN INTERVAL LENGTH IS 2*(-MAXSTK)*(B-A). ARRAYS ADA 950
C          XLEFT XRIGHT                             ADA 960
C          ADA 970
C          ***** PORTABILITY CONSIDERATIONS ***** ADA 980
C          ADA 990
C          THIS PROGRAM IS IN ANSI STANDARD FORTRAN. IN ADDITION, IT MEETS ADA 1000
C          ALL THE REQUIREMENTS OF THE BELL LABS PORTABLE FORTRAN -PFORT- ADA 1010
C          EXCEPT ONE. HOLLERITH CHARACTERS ARE PACKED 4/WORD RATHER THAN ADA 1020
C          1/WORD AS SPECIFIED BY PFORT.             ADA 1030
C          NEVERTHELESS, THIS PROGRAM IS AFFECTED IN SEVERAL WAYS BY A ADA 1040
C          CHANGE IN MACHINE WORD LENGTH AND CHANGING TO DOUBLE PRECISION. ADA 1050
C          ***** THIS VERSION IS FOR THE MACHINE WITH THE LONGEST SINGLE ADA 1060
C          PRECISION WORD (CDC). THE LENGTH OF SOME CONSTANTS IN ADA 1070
C          THE SUBPROGRAM COMPUT EXCEEDS THE CAPACITY OF SOME FORTRAN ADA 1080
C          COMPILERS AND CAN PREVENT COMPILEATION. ADA 1090
C          INPUT-OUTPUT -- IS OF THREE TYPES.        ADA 1100
C          FATAL ERROR MESSAGES - OCCUR IN SETUP,TAKE,PUT AND TERMIN ADA 1110
C          THEY CANNOT BE SWITCHED OFF              ADA 1120
C          USER AND DEBUGGING OUTPUT - CAN BE SWITCHED OFF ADA 1130
C          THESE INVOLVE MANY FORMATS LIKE E15.8, F12.8, E22.13, ETC. ADA 1140
C          SOME FORTRAN COMPILERS REQUIRE D-FORMAT FOR DOUBLE PRECISION. ADA 1150
C          SOME DO NOT HANDLE E22.13 ON MACHINES OF SHORTER WORD LENGTH. ADA 1160
C          ADA 1170
C          SUMARY PRINTS COEFFICIENTS 5 PER LINE, 6 PER LINE IS BETTER ADA 1180
C          FOR SHORTER WORD LENGTHS. DOUBLE PRECISION ON MANY MACHINES ADA 1190
C          LIMITS ONE TO 4 PER LINE.                ADA 1200
C          ADA 1210
C          CONSTANTS -- THE GAUSS WEIGHTS AND ABSCISSA IN COMPUTE ARE GIVEN ADA 1220
C          TO 15 DIGITS IN THE COMMENTS             ADA 1230
C          ADA 1240
C          DOUBLE PRECISION CONVERSION -- REQUIRES FOUR STEPS ADA 1250
C          1. ALL REAL VARIABLES ARE DECLARED DOUBLE PRECISION. THIS IS ADA 1260
C          DONE BY CHANGING REAL TO DOUBLE PRECISION AS ALL REALS ARE ADA 1270
C          EXPLICITLY DECLARED AND ROOM IS LEFT FOR THIS CHANGE. REAL ADA 1280
C          VARIABLES APPEAR BEFORE INTEGERS IN ALL COMMON BLOCKS. ADA 1290
C          ADA 1300
C          2. ADD D TO CONSTANTS( E.G. 1.0 = 1.0D0 ). ADJUST LENGTH OF ADA 1310
C          GAUSS WEIGHTS IN COMPUTE.                 ADA 1320
C          ADA 1330
C          3. CHANGE ABS,AMAX1,FLOAT AT MANY PLACES ADA 1340
C          ADA 1350

```



```

C      4. ADJUST THE INTERVAL STACK SIZE = DIMENSIONS OF XLEFT, XRIGHT ADA 1360
C      AND VALUE OF MAXSTK. ADJUST THE VALUE BUFFER IN PUT TO BE ADA 1370
C      ABOUT .1E-K FOR A MACHINE WITH K+2 DECIMAL DIGITS. ADA 1380
C      ADA 1390
C      REAL A, ACCUR, B, BLEFT, BRIGHT, CHARF, DDTEMP, DSCTOL, ERROR, ADA 1400
*      ERROR, FACTOR, FINTRP, FLEFT, FRIGHT, NORM, XBREAK, XDD, ADA 1410
*      XINTRP, XLEFT, XRIGHT ADA 1420
      DIMENSION XBREAK(20), DBREAK(20), BLEFT(20), BRIGHT(20) ADA 1430
      DIMENSION XLEFT(50), XRIGHT(50), FACTOR(12), FMESGE(6) ADA 1440
      DIMENSION DDTEMP(12,12), FINTRP(10), FLEFT(6), FRIGHT(6), ADA 1450
*      XDD(12), XINTRP(10) ADA 1460
      INTEGER BOTH, BREAK, DBREAK, DEGREE, EDIST, FMESGE, RIGHT, ADA 1470
*      RIGHTX, SMOOTH ADA 1480
      LOGICAL DISCRD, FATAL, FINISH ADA 1490
      REAL XKNOTS(KDIMEN), COEFS(KDIMEN,NDIMEN) ADA 1500
      EXTERNAL F ADA 1510
      REAL F ADA 1520
C      ADA 1530
C      COMMON /INPUTZ/ A, B, ACCUR, NORM, CHARF, XBREAK, BLEFT, BRIGHT, ADA 1540
*      DBREAK, DEGREE, SMOOTH, LEVEL, EDIST, NBREAK, KNTDIM, NPARDM ADA 1550
C      KNTDIM - KDIMEN, NAME CHANGED TO PUT IN COMMON ADA 1560
C      NPARDM - NDIMEN, NAME CHANGED TO PUT IN COMMON ADA 1570
C      COMMON /RESULZ/ ERROR, KNOTS ADA 1580
C      KNOTS = FINAL NO. OF KNOTS, INCLUDES B AS ONE. ADA 1590
C      ERROR = ESTIMATE OF ERROR ACTUALLY ACHIEVED. ADA 1600
C      COMMON /KONTRL/ DSCTOL, ERRORI, XLEFT, XRIGHT, BREAK, BOTH, ADA 1610
*      FACTOR, FMESGE, IBREAK, INTERP, LEFT, MAXAUX, MAXKNT, MAXPAR, ADA 1620
*      MAXSTK, NPAR, NSTACK, RIGHT, DISCRD, FATAL, FINISH ADA 1630
C      KONTRL CONTAINS GENERALLY USEFUL VARIABLES ADA 1640
C      FATAL - SWITCH FOR DETECTION OF FATAL ERROR ADA 1650
C      INCLUDING EXCESSIVE INTERVAL SUBDIVISION ADA 1660
C      WHICH DOES NOT TERMINATE THE COMPUTATION ADA 1670
C      FINISH - SWITCH TO TERMINATE ALGORITHM ADA 1680
C      MAXS - SEE COMMENTS EARLIER ADA 1690
C      NSTACK - COUNTER FOR INTERVAL STACK, CONSISTS OF ADA 1700
C      (XLEFT(J),XRIGHT(J)) J = 1 TO NSTACK ADA 1710
C      ERRORI - ERROR ESTIMATE FOR TOP INTERVAL ADA 1720
C      DSCTOL - TOLERANCE TO CHECK DISCARDING INTERVALS ADA 1730
C      DISCRD - SWITCH TO SIGNAL DISCARD OF TOP INTERVAL ADA 1740
C      FACTOR - ARRAY OF FACTORIALS ADA 1750
C      NPAR - NUMBER OF PAREMETERS = DEGREE + 1 ADA 1760
C      FMESGE - STRING = **** FATAL ERROR **** ADA 1770
C      INTERP - NUMBER OF INTERIOR INTERPOLATION POINTS ADA 1780
C      IN THE NORMAL INTERVAL ADA 1790
C      IBREAK - COUNTER ON BREAK POINTS ADA 1800
C      BREAK - SWITCH FOR BREAK POINT IN TOP INTERVAL ADA 1810
C      0 = NO BREAK PRESENT ADA 1820
C      LEFT = BREAK AT XLEFT(NSTACK) ADA 1830
C      RIGHT = BREAK AT XRIGHT(NSTACK) ADA 1840
C      BOTH = BREAK AT BOTH ENDS ADA 1850
C      COMMON /COMDIF/ DDTEMP, FINTRP, FLEFT, FRIGHT, XDD, XINTRP, ADA 1860
*      LEFTX, NINTRP, RIGHTX ADA 1870
C      COMDIF CONTAINS VARIABLES USED ONLY BY COMPUT AND FRIENDS. ADA 1880
C      NINTRP - NUMBER OF INTERIOR INTERPOLATION POINTS ADA 1890
C      FOR THE CURRENT INTERVAL ADA 1900
C      XINTRP - INTERIOR INTERPOLATION POINTS ADA 1910
C      FINTRP - F VALUES AT XINTRP POINTS ADA 1920
C      LEFTX - MULTIPLICITY OF INTERPOLATION AT XLEFT ADA 1930
C      = NO. OF DERIVATIVES MATCHED AT XLEFT ADA 1940
C      FLEFT - VALUES OF F AND ITS DERIVATIVES AT XLEFT ADA 1950
C      RIGHTX - MULTIPLICITY OF INTERPOLATION AT XRIGHT ADA 1960
C      FRIGHT - VALUES OF F AND DERIVATIVES AT XRIGHT ADA 1970
C      DDTEMP - THE ARRAY OF DIVIDED DIFFERENCES ADA 1980
C      XDD - THE X VALUES FOR DDTEMP WITH PROPER ADA 1990
C      MULTIPLICITIES OF XLEFT AND XRIGHT ADA 2000
C      COMMON /SAVEIT/ IKNOT ADA 2010
C      ADA 2020
C----- MAIN CONTROL PROGRAM ----- ADA 2030
C      ADA 2040
C      ***** NOTE - ARGUMENTS BELOW ARE FOR READABILITY ONLY ***** ADA 2050
C      ***** EXCEPT FOR F AND XKNOTS,COEFS,KDIMEN,NDIMEN ***** ADA 2060
C      ADA 2070
C      CHECK INPUT, INITIAL COMPUTATIONS, PRINT PROBLEM ADA 2080
C      ADA 2090
C      CALL SETUP(XKNOTS, COEFS, KDIMEN, NDIMEN) ADA 2100
C      ADA 2110
C      CHECK FOR FATAL ERROR IN PROBLEM SPECIFICATION ADA 2120
C      IF (FATAL) RETURN ADA 2130

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C
C          LOOP OVER PROCESSING OF INTERVALS
C 10 CALL TAKE(INTERV)
C
C          CALL COMPUT(F, APPROX, INTERV)
C
C          CHECK FOR DISCARDING INTERVALS
C CALL CHECK(FUNCT, CHARCT)
C
C          PUT NEW INTERVALS ON STACK OR DISCARD, UPDATE STATUS
C CALL PUT(INTERV, XKNOTS, COEFS, KDIMEN, NDIMEN)
C
C CALL TERMIN(TEST, AND, PRINT, XKNOTS, KDIMEN)
C
C IF (.NOT.FINISH) GO TO 10
C
C CALL SUMARY(XKNOTS, COEFS, KDIMEN, NDIMEN)
C RETURN
C END
C
C          SUBROUTINE PPFIT4(F, XLFT, XRGT, EPSLN, NPIECE, ERREST, XKNOTS,
* COEFS, KDIMEN, NDIMEN, NDEG, NSMTH, EMEAS, LPRNT, FOSCIL, ATYPE,
* KBREAK, BRAKPT, KDERVB, VALLFT, VALRGT)
C
C          =====PPF 50
C          PPF 60
C ** THIS SET OF FOUR CONTROL PROGRAMS SET VARYING NUMBERS OF DEFAULT PPF 70
C VALUES FOR ARGUMENTS. IT USES ENTRY STATEMENTS WHICH ARE DONE PPF 80
C DIFFERENTLY BY VARIOUS FORTRANS. FOR THIS REASON ENTRY STATEMENTS PPF 90
C ARE ONLY INDICATED BY COMMENT CARDS. WRITING FOUR SEPARATE ROU- PPF 100
C TINES APPROXIMATELY TRIPLES THE LENGTH OF THE TOTAL CODE. PPF 110
C PPF 120
C THE FOLLOWING TABULATES THE INTERNAL AND EXTERNAL NAMES OF THE PPF 130
C ARGUMENTS ALONG WITH THEIR DEFAULT VALUES FOR THE VARIOUS PPFIT. PPF 140
C NOTE THAT THIS ALLOWS THE ARGUMENTS TO BE PUT INTO A COMMON BLOCK PPF 150
C AND AVOIDS LONG ARGUMENT LISTS WITHIN ADAPT ITSELF. PPF 160
C PPF 170
C PPF 180
C INTERNAL   EXTERNAL   DEFAULT VALUE   SET BY PPFIT NUMBER   PPF 190
C F          F          NONE                PPF 200
C A,B       A,B       NONE                PPF 210
C ACCUR     EPSLN     NONE                PPF 220
C KNOTS     NPIECE    OUTPUT             PPF 230
C ERROR     ERREST    OUTPUT             PPF 240
C XKNOTS    XKNOTS    OUTPUT             PPF 250
C COEFS     COEFS     OUTPUT             PPF 260
C KDIMEN    KDIMEN    NONE                PPF 270
C NDIMEN    NDIMEN    NONE                PPF 280
C DEGREE    NDEG      3                  1                    PPF 290
C SMOOTH    NSMTH     0                  1                    PPF 300
C NORM      EMEAS     3                  1 2                  PPF 310
C LEVEL     LPRNT     1                  1 2                  PPF 320
C CHARF     FOSCIL    VARIABLE           1 2                  PPF 330
C EDIST     ATYPE     1                  1 2 3                PPF 340
C NBREAK    KBREAK    0                  1 2 3                PPF 350
C XBREAK    BRAKPT    -                   PPF 360
C DBREAK    KDERVB    -                   PPF 370
C BLEFT     VALLFT    -                   PPF 380
C BRIGHT    VALRGT    -                   PPF 390
C PPF 400
C REAL A, ACCUR, B, BLEFT, BRIGHT, CHARF, DDTEMP, DSCTOL, ERROR, PPF 410
* ERRORI, FACTOR, FINTRP, FLEFT, FRIGHT, NORM, XBREAK, XDD, PPF 420
* XINTRP, XLEFT, XRIGHT PPF 430
C DIMENSION XBREAK(20), DBREAK(20), BLEFT(20), BRIGHT(20) PPF 440
C DIMENSION XLEFT(50), XRIGHT(50), FACTOR(12), FMESGE(6) PPF 450
C DIMENSION DDTEMP(12,12), FINTRP(10), FLEFT(6), FRIGHT(6), PPF 460
* XDD(12), XINTRP(10) PPF 470
C INTEGER BOTH, BREAK, DBREAK, DEGREE, EDIST, FMESGE, RIGHT, PPF 480
* RIGHTX, SMOOTH PPF 490
C LOGICAL DISCRD, FATAL, FINISH PPF 500
C REAL XKNOTS(KDIMEN), COEFS(KDIMEN,NDIMEN) PPF 510
C REAL BRAKPT, EMEAS, EPSLN, ERREST, FOSCIL, F, VALLFT, VALRGT, PPF 520
* XLFT, XRGT PPF 530
C DIMENSION BRAKPT(20), KDERVB(20), VALLFT(20), VALRGT(20) PPF 540
C INTEGER ATYPE PPF 550
C EXTERNAL F PPF 560

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COMMON /INPUTZ/ A, B, ACCUR, NORM, CHARF, XBREAK, BLEFT, BRIGHT, PPF 570
* DBREAK, DEGREE, SMOOTH, LEVEL, EDIST, NBREAK, KNTDIM, NPARDM PPF 580
COMMON /RESULZ/ ERROR, KNOTS PPF 590
COMMON /KONTRL/ DSCTOL, ERRORI, XLEFT, XRIGHT, BREAK, BOTH, PPF 600
* FACTOR, FMESGE, IBREAK, INTERP, LEFT, MAXAUX, MAXKNT, MAXPAR, PPF 610
* MAXSTK, NPAR, NSTACK, RIGHT, DISCRD, FATAL, FINISH PPF 620
COMMON /COMDIF/ DDTEMP, FINTRP, FLEFT, FRIGHT, XDD, XINTRP, PPF 630
* LEFTX, NINTRP, RIGHTX PPF 640
COMMON /SAVEIT/ IKNOT PPF 650
C PPF 660
C PPF 670
C          SKIP ALL DEFAULT SETTING FOR PPFIT4 PPF 680
GO TO 10 PPF 690
C PPF 700
C          SET DEFAULTS FOR PPFIT1 PPF 710
C PPF 720
C          ***** SINCE ENTRY IS NON-STANDARD ***** PPF 730
C          ***** THE NATURAL WAY TO IMPLEMENT ***** PPF 740
C          ***** THE OTHER PPFIT5 IS ONLY ***** PPF 750
C          ***** INDICATED IN THE COMMENTS ***** PPF 760
C PPF 770
C          ENTRY PPFIT1 PPF 780
C          ENTRY PPFIT1(F,XLFT,XRGT,EPSLN,NPIECE,ERREST,XKNOTS,COEFS, PPF 790
C          A          KDIMEN,NDIMEN) PPF 800
          NDEG = 3 PPF 810
          NSMTH = 0 PPF 820
C PPF 830
C          SET DEFAULTS FOR PPFIT2 PPF 840
C          ENTRY PPFIT2 PPF 850
C          ENTRY PPFIT2(F,XLFT,XRGT,EPSLN,NPIECE,ERREST,XKNOTS,COEFS, PPF 860
C          A          KDIMEN,NDIMEN,NDEG,NSMTH) PPF 870
          EMEAS = 3.0 PPF 880
          LPRNT = 1 PPF 890
          FOSCIL = B - A PPF 900
          IF (NDEG.EQ.2) FOSCIL = .5*FOSCIL PPF 910
          IF (NDEG.LE.1) FOSCIL = (B-A)/3. PPF 920
C PPF 930
C          SET DEFAULTS FOR PPFIT3 PPF 940
C          ENTRY PPFIT3 PPF 950
C          ENTRY PPFIT3(F,XLFT,XRGT,EPSLN,NPIECE,ERREST,XKNOTS,COEFS, PPF 960
C          A          KDIMEN,NDIMEN,NDEG,NSMTH,EMEAS,LPRNT,FOSCIL) PPF 970
          ATYPE = 1 PPF 980
          KSING = 0 PPF 990
          KBREAK = 0 PPF 1000
C PPF 1010
C          PUT INPUT INTO COMMON INPUTZ BY CHANGING TO INTERNAL NAMES PPF 1020
10 A = XLFT PPF 1030
   B = XRGT PPF 1040
   ACCUR = EPSLN PPF 1050
   DEGREE = NDEG PPF 1060
   SMOOTH = NSMTH PPF 1070
   NORM = EMEAS PPF 1080
   LEVEL = LPRNT PPF 1090
   CHARF = FOSCIL PPF 1100
   EDIST = ATYPE PPF 1110
   NBREAK = KBREAK PPF 1120
   IF (NBREAK.LE.0 .OR. NBREAK.GE.21) GO TO 30 PPF 1130
   DO 20 K=1,NBREAK PPF 1140
     XBREAK(K) = BRAKPT(K) PPF 1150
     DBREAK(K) = KDERVB(K) PPF 1160
     BLEFT(K) = VALLFT(K) PPF 1170
     BRIGHT(K) = VALRGT(K) PPF 1180
20 CONTINUE PPF 1190
30 CONTINUE PPF 1200
C PPF 1210
C          CALL ADAPT(F, XKNOTS, COEFS, KDIMEN, NDIMEN) PPF 1220
          NPIECE = KNOTS PPF 1230
          ERREST = ERROR PPF 1240
          RETURN PPF 1250
          END PPF 1260

SUBROUTINE CHECK(FUNCT, CHAR) CHE 10
C CHE 20

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C -----CHE 30
C CHE 40
C ** THIS PROGRAM CHECKS FOR DISCARDING INTERVAL, APPLIES VARIOUS CHE 50
C TESTS ABOUT DISCARDING INVOLVING EDIST AND CHARF. CHE 60
C CHE 70
C REAL A, ACCUR, B, BLEFT, BRIGHT, CHARF, DDTEMP, DSCTOL, ERROR, CHE 80
* ERRORI, FACTOR, FINTRP, FLEFT, FRIGHT, NORM, XBREAK, XDD, CHE 90
* XINTRP, XLEFT, XRIGHT CHE 100
DIMENSION XBREAK(20), DBREAK(20), BLEFT(20), BRIGHT(20) CHE 110
DIMENSION XLEFT(50), XRIGHT(50), FACTOR(12), FMESGE(6) CHE 120
DIMENSION DDTEMP(12,12), FINTRP(10), FLEFT(6), FRIGHT(6), CHE 130
* XDD(12), XINTRP(10) CHE 140
INTEGER BOTH, BREAK, DBREAK, DEGREE, EDIST, FMESGE, RIGHT, CHE 150
* RIGHTX, SMOOTH CHE 160
LOGICAL DISCRD, FATAL, FINISH CHE 170
REAL DTEST, DX CHE 180
COMMON /INPUTZ/ A, B, ACCUR, NORM, CHARF, XBREAK, BLEFT, BRIGHT, CHE 190
* DBREAK, DEGREE, SMOOTH, LEVEL, EDIST, NBREAK, KNTDIM, NPARDM CHE 200
COMMON /RESULZ/ ERROR, KNOTS CHE 210
COMMON /KONTRL/ DSCTOL, ERRORI, XLEFT, XRIGHT, BREAK, BOTH, CHE 220
* FACTOR, FMESGE, IBREAK, INTERP, LEFT, MAXAUX, MAXKNT, MAXPAR, CHE 230
* MAXSTK, NPAR, NSTACK, RIGHT, DISCRD, FATAL, FINISH CHE 240
COMMON /COMDIF/ DDTEMP, FINTRP, FLEFT, FRIGHT, XDD, XINTRP, CHE 250
* LEFTX, NINTRP, RIGHTX CHE 260
C CHE 270
C CHE 280
C DISCRD = .FALSE. CHE 290
DX = XRIGHT(NSTACK) - XLEFT(NSTACK) CHE 300
C CHE 310
C COMPUTE DTEST FOR IMPLEMENTING VARIOUS TYPES OF ADAPTIVE APPROX CHE 320
IF (EDIST.EQ.0) DTEST = DX*DSCTOL CHE 330
C FOR THE APPROXIMATE FIXED ERROR DISTRIBUTION TYPE WE ESTIMATE CHE 340
C THE FINAL NUMBER OF KNOTS BY( LIMITING IT A LITTLE ) CHE 350
C (NSTACK+KNOTS+2)((B-A)/(XRIGHT-A)) CHE 360
IF (EDIST.EQ.1) DTEST = DSCTOL/(FLOAT(NSTACK+KNOTS+2)*AMINI((B-A)/CHE 370
* (XRIGHT(NSTACK)-A),5.)) CHE 380
IF (EDIST.EQ.2 .OR. NORM.EQ.3.) DTEST = DSCTOL CHE 390
C CHE 400
C CHECK FOR DISCARD OF INTERVAL CHE 410
IF (ERRORI.LE.DTEST) DISCRD = .TRUE. CHE 420
C CHE 430
C CHECK CHARACTERISTIC LENGTH OF FUNCTION CHE 440
IF (DX.GE.CHARF) DISCRD = .FALSE. CHE 450
RETURN CHE 460
END CHE 470

SUBROUTINE COMPUT(F, APPROX, INTERV) COM 10
C COM 20
C -----COM 30
C COM 40
C ** THIS PROGRAM COMPUTES THE PIECEWISE POLYNOMIAL APPROXIMATION ON COM 50
C THE CURRENT INTERVAL. IT ALSO ESTIMATES THE ERROR COM 60
C COM 70
C REAL A, ACCUR, B, BLEFT, BRIGHT, CHARF, DDTEMP, DSCTOL, ERROR, COM 80
* ERRORI, FACTOR, FINTRP, FLEFT, FRIGHT, NORM, XBREAK, XDD, COM 90
* XINTRP, XLEFT, XRIGHT COM 100
DIMENSION XBREAK(20), DBREAK(20), BLEFT(20), BRIGHT(20) COM 110
DIMENSION XLEFT(50), XRIGHT(50), FACTOR(12), FMESGE(6) COM 120
DIMENSION DDTEMP(12,12), FINTRP(10), FLEFT(6), FRIGHT(6), COM 130
* XDD(12), XINTRP(10) COM 140
INTEGER BOTH, BREAK, DBREAK, DEGREE, EDIST, FMESGE, RIGHT, COM 150
* RIGHTX, SMOOTH COM 160
LOGICAL DISCRD, FATAL, FINISH COM 170
REAL ABSC, DX, ERRINT, F, FDERVL, FDERVR, FDUMB, POLYDD, WGT5 COM 180
DIMENSION ABSC(4), WGT5(4), FDERVL(5), FDERVR(5), FDUMB(5) COM 190
EXTERNAL F, POLYDD COM 200
COMMON /INPUTZ/ A, B, ACCUR, NORM, CHARF, XBREAK, BLEFT, BRIGHT, COM 210
* DBREAK, DEGREE, SMOOTH, LEVEL, EDIST, NBREAK, KNTDIM, NPARDM COM 220
COMMON /RESULZ/ ERROR, KNOTS COM 230
COMMON /KONTRL/ DSCTOL, ERRORI, XLEFT, XRIGHT, BREAK, BOTH, COM 240
* FACTOR, FMESGE, IBREAK, INTERP, LEFT, MAXAUX, MAXKNT, MAXPAR, COM 250
* MAXSTK, NPAR, NSTACK, RIGHT, DISCRD, FATAL, FINISH COM 260
COMMON /COMDIF/ DDTEMP, FINTRP, FLEFT, FRIGHT, XDD, XINTRP, COM 270
* LEFTX, NINTRP, RIGHTX COM 280

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C                                     COM 290
      EQUIVALENCE (FLEFT(2),FDERVL(1)), (FRIGHT(2),FDERVR(1))      COM 300
C      FIFTEEN DIGIT VALUES FOR THESE GAUSS INTEGRATION CONSTANTS ARE      COM 310
C      .861136311594053 .339981043584856 .347854845137454 .652145154862546      COM 320
      DATA ABSC(1) /-.86113631159405/      COM 330
      DATA ABSC(2) /-.33998104358486/      COM 340
      DATA ABSC(3) /.33998104358486/      COM 350
      DATA ABSC(4) /.86113631159405/      COM 360
      DATA WGTS(1) /.34785484513745/      COM 370
      DATA WGTS(2) /.65214515486255/      COM 380
      DATA WGTS(3) /.65214515486255/      COM 390
      DATA WGTS(4) /.34785484513745/      COM 400
C                                     COM 410
C      COMPUTE INTERPOLATION INFORMATION      COM 420
      NINTRP = DEGREE - 2*SMOOTH - 1      COM 430
C                                     COM 440
C      INCREASE NUMBER OF INTERPOLATION POINTS IF BREAK POINTS ARE      COM 450
C      SPECIFIED WITH FEWER DERIVATIVES THAN SMOOTH      COM 460
      IF (BREAK.EQ.LEFT .OR. BREAK.EQ.RIGHT) NINTRP = NINTRP + SMOOTH -      COM 470
      * DBREAK(IBREAK)      COM 480
      IF (BREAK.EQ.BOTH) NINTRP = NINTRP + 2*SMOOTH - DBREAK(IBREAK) -      COM 490
      * DBREAK(IBREAK+1)      COM 500
      IF (NINTRP.EQ.0) GO TO 20      COM 510
C      GENERATE EQUAL SPACED INTERPOLATION POINTS      COM 520
      DX = (XRIGHT(NSTACK)-XLEFT(NSTACK))/FLOAT(NINTRP+1)      COM 530
      DO 10 J=1,NINTRP      COM 540
          XINTRP(J) = XLEFT(NSTACK) + FLOAT(J)*DX      COM 550
10 CONTINUE      COM 560
C                                     COM 570
C      GET LEFT AND RIGHT F-VALUES, PUT F-VALUE IN FIRST ELEMENT      COM 580
C      OF ARRAYS FLEFT AND FRIGHT. GET DERIVATIVES BACK AS      COM 590
C      OTHER ELEMENTS VIA THE SUBARRAYS FDERVL AND FDERVR.      COM 600
20 FLEFT(1) = F(XLEFT(NSTACK),FDERVL)      COM 610
   FRIGHT(1) = F(XRIGHT(NSTACK),FDERVR)      COM 620
   LEFTX = SMOOTH + 1      COM 630
   RIGHTX = LEFTX      COM 640
C      GET F-VALUES AT OTHER INTERPOLATION POINTS, IF ANY      COM 650
      IF (NINTRP.EQ.0) GO TO 40      COM 660
      DO 30 J=1,NINTRP      COM 670
          FINTRP(J) = F(XINTRP(J),FDUMB)      COM 680
30 CONTINUE      COM 690
C                                     COM 700
C      CHECK FOR BREAK POINTS, MODIFY VALUES IF NECESSARY      COM 710
40 CONTINUE      COM 720
      IF (BREAK.NE.LEFT) GO TO 50      COM 730
      LEFTX = DBREAK(IBREAK) + 1      COM 740
      FLEFT(LEFTX) = FRIGHT(IBREAK)      COM 750
50 IF (BREAK.NE.RIGHT) GO TO 60      COM 760
      RIGHTX = DBREAK(IBREAK) + 1      COM 770
      FRIGHT(RIGHTX) = FLEFT(IBREAK)      COM 780
60 IF (BREAK.NE.BOTH) GO TO 70      COM 790
      LEFTX = DBREAK(IBREAK) + 1      COM 800
      RIGHTX = DBREAK(IBREAK+1) + 1      COM 810
      FLEFT(LEFTX) = FRIGHT(IBREAK)      COM 820
      FRIGHT(RIGHTX) = FLEFT(IBREAK+1)      COM 830
70 CONTINUE      COM 840
C                                     COM 850
C      COMPUTE DIVIDED DIFFERENCES, NEWTON FORM OF POLYNOMIAL      COM 860
      CALL NEWTON(LEFTX, RIGHTX, NINTRP)      COM 870
C                                     COM 880
C      COMPUTE NORM OF ERROR OF THIS APPROXIMATION USING FOUR PTS      COM 890
C      ADD 50 PERCENT FUDGE FACTOR      COM 900
      ERRORI = ERRINT(F, POLYDD, XLEFT(NSTACK), XRIGHT(NSTACK), ABSC, WGTS)      COM 910
      ERRORI = 1.5*ERRORI      COM 920
      RETURN      COM 930
      END      COM 940

REAL FUNCTION ERRINT(F, FIT, AAA, BBB, POINTS, WEIGHT)      ERR 10
C                                     ERR 20
C      -----ERR 30
C                                     ERR 40
C ** THIS FUNCTION DOES A FOUR POINT INTEGRATION RULE FOR THE      ERR 50
C ABSOLUTE VALUE OF THE DIFFERENCE OF TWO FUNCTIONS( F AND FIT )      ERR 60
C      ABS( F(X) - FIT(X) )**NORM      ERR 70
C      THE INTEGRATION USES THE POINTS AND WEIGHTS GIVEN AND SCALED      ERR 80

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C      FROM (-1,1) TO (AAA,BBB)                                ERR  90
C                                                                 ERR 100

      SUBROUTINE NEWTON(NL, NR, NI)                             NEW  10
C                                                                 NEW  20
C      =====NEW  30
C                                                                 NEW  40
C ** THIS PROGRAM COMPUTES THE DIVIDED DIFFERENCES ARRAY AS FOLLOWS NEW  50
C      NL COALESCED POINTS ON LEFT - DERIV VALUES IN FLEFT NEW  60
C      NR COALESCED POINTS ON RIGHT - - - - - FRIGHT NEW  70
C      NI DISTINCT POINTS INBETWEEN - FNCTN - - - FINTRP NEW  80
C                                                                 NEW  90
C      THE POINTS ARE ORDERED XL = XLEFT (NSTACK) NEW 100
C                                                                 XR = XRIGHT(NSTACK) NEW 110
C                                                                 XINTRP ARRAY NEW 120
C                                                                 NEW 130
C      LAYOUT OF THE DDTEMP DIVIDED DIFFERENCE ARRAY NEW 140
C                                                                 NEW 150
C      NL=6  LLLLLL****II NEW 160
C      NR=4  LLLLL****II   L = FIRST TRIANGLE NEW 170
C      NI=2  LLL****II NEW 180
C           LLL****II   R = SECOND TRIANGLE NEW 190
C           LL****II NEW 200
C           L****II   * = FILL BETWEEN TRIANGLES NEW 210
C           RRRRII NEW 220
C           RRRRII   I = COMPLETION FOR INTERPOLATION POINTS NEW 230
C           RRII NEW 240
C           RII   IDIF = HORIZONTAL COORD. = DIFFERENCE ORDER NEW 250
C           II   IPT = VERTICAL COORD. ASSOCIATED WITH POINTS NEW 260
C           I NEW 270
C                                                                 NEW 280
C                                                                 NEW 290
C      REAL A, ACCUR, B, BLEFT, BRIGHT, CHARF, DDTEMP, DSCTOL, ERROR, NEW 300
C      * ERRORI, FACTOR, FINTRP, FLEFT, FRIGHT, NORM, XBREAK, XDD, NEW 310
C      * XINTRP, XLEFT, XRIGHT NEW 320
C      DIMENSION XBREAK(20), DBREAK(20), BLEFT(20), BRIGHT(20) NEW 330
C      DIMENSION XLEFT(50), XRIGHT(50), FACTOR(12), FMESGE(6) NEW 340
C      DIMENSION DDTEMP(12,12), FINTRP(10), FLEFT(6), FRIGHT(6), NEW 350
C      * XDD(12), XINTRP(10) NEW 360
C      INTEGER BOTH, BREAK, DBREAK, DEGREE, EDIST, FMESGE, RIGHT, NEW 370
C      * RIGHTX, SMOOTH NEW 380
C      LOGICAL DISCRD, FATAL, FINISH NEW 390
C      REAL DIFF, DIFFX NEW 400
C      COMMON /INPUTZ/ A, B, ACCUR, NORM, CHARF, XBREAK, BLEFT, BRIGHT, NEW 410
C      * DBREAK, DEGREE, SMOOTH, LEVEL, EDIST, NBREAK, KNTDIM, NPARDM NEW 420
C      COMMON /RESULZ/ ERROR, KNOTS NEW 430
C      COMMON /KONTRL/ DSCTOL, ERRORI, XLEFT, XRIGHT, BREAK, BOTH, NEW 440
C      * FACTOR, FMESGE, IBREAK, INTERP, LEFT, MAXAUX, MAXKNT, MAXPAR, NEW 450
C      * MAXSTK, NPAR, NSTACK, RIGHT, DISCRD, FATAL, FINISH NEW 460
C      COMMON /COMDIF/ DDTEMP, FINTRP, FLEFT, FRIGHT, XDD, XINTRP, NEW 470
C      * LEFTX, NINTRP, RIGHTX NEW 480
C                                                                 NEW 490
C      MAIN CALCULATION OF DIVIDED DIFFERENCES NEW 500
C      DEFINE A FEW SHORT CONSTANTS NEW 510
C      NL1 = NL - 1 NEW 520
C      NL2 = NL + 1 NEW 530
C      NR1 = NR - 1 NEW 540
C      NR2 = NR + 1 NEW 550
C      NRL = NR + NL NEW 560
C                                                                 NEW 570
C      PUT X-VALUES IN A SINGLE ARRAY WITH NDDX = NL+NR+NI POINTS NEW 580
C      DO 10 NDDX=1,NL NEW 590
C          XDD(NDDX) = XLEFT(NSTACK) NEW 600
C 10 CONTINUE NEW 610
C          NDDX = NL NEW 620
C          DO 20 K=1,NR NEW 630
C              NDDX = NDDX + 1 NEW 640
C              XDD(NDDX) = XRIGHT(NSTACK) NEW 650
C 20 CONTINUE NEW 660
C      CHECK IF THERE ARE ANY INTERPOLATION POINTS TO ADD TO XDD NEW 670
C      IF (NI.EQ.0) GO TO 40 NEW 680
C      DO 30 K=1,NI NEW 690
C          NDDX = NDDX + 1 NEW 700
C          XDD(NDDX) = XINTRP(K) NEW 710
C 30 CONTINUE NEW 720

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C		NEW	730
C	FILL BORDER OF FIRST TRIANGLE - SIZE NL.	NEW	740
40	CONTINUE	NEW	750
C	TOP BORDER	NEW	760
	DO 50 IDIF=1,NL	NEW	770
	DDTEMP(IDIF,1) = FLEFT(IDIF)/FACTOR(IDIF)	NEW	780
50	CONTINUE	NEW	790
	IF (NL1.EQ.0) GO TO 70	NEW	800
C	BOTTOM BORDER	NEW	810
	DO 60 IDIF=1,NL1	NEW	820
	IPT = NL2 - IDIF	NEW	830
	DDTEMP(IDIF,IPT) = DDTEMP(IDIF,1)	NEW	840
60	CONTINUE	NEW	850
		NEW	860
C	FILL BORDER OF SECOND TRIANGLE - SIZE NR	NEW	870
70	CONTINUE	NEW	880
C	TOP BORDER	NEW	890
	DO 80 IDIF=1,NR	NEW	900
	DDTEMP(IDIF,NL2) = FRIGHT(IDIF)/FACTOR(IDIF)	NEW	910
80	CONTINUE	NEW	920
	IF (NRL.EQ.NL2) GO TO 100	NEW	930
C	BOTTOM BORDER	NEW	940
	DO 90 IDIF=1,NR1	NEW	950
	IPT = NRL + 1 - IDIF	NEW	960
	DDTEMP(IDIF,IPT) = DDTEMP(IDIF,NL2)	NEW	970
90	CONTINUE	NEW	980
		NEW	990
C	FILL PARALLOGRAM BETWEEN THE TWO TRIANGLES JUST FILLED	NEW	1000
C	FILL ENTRIES PARALLEL TO BOTTOM OF FIRST TRIANGLE	NEW	1010
100	CONTINUE	NEW	1020
C		NEW	1030
C	LOOP STEPPING ALONG TOP SIDE OF SECOND TRIANGLE	NEW	1040
	DO 120 J=2,NR2	NEW	1050
	IDIF = J	NEW	1060
C	LOOP STEPPING PARALLEL TO BOTTOM SIDE OF FIRST TRIANGLE	NEW	1070
	DO 110 K=2,NL2	NEW	1080
	IPT = NL + 2 - K	NEW	1090
	DIFFF = DDTEMP(IDIF-1,IPT+1) - DDTEMP(IDIF-1,IPT)	NEW	1100
	IPT2 = IPT - 1 + IDIF	NEW	1110
	DIFFX = XDD(IPT2) - XDD(IPT)	NEW	1120
	DDTEMP(IDIF,IPT) = DIFFF/DIFFX	NEW	1130
	IDIF = IDIF + 1	NEW	1140
110	CONTINUE	NEW	1150
120	CONTINUE	NEW	1160
		NEW	1170
C	DEBUG DEBUG DEBUG DEBUG	NEW	1170
	IF (LEVEL.GE.4 .AND. KNOTS.LE.1) WRITE (6,99999) NR2, NL2, IDIF,	NEW	1180
	* IPT, DIFFF, DIFFX	NEW	1190
		NEW	1200
C	FILL IN BOTTOM DIAGONALS FOR INTERPOLATION POINTS, IF ANY	NEW	1210
	IF (NI.EQ.0) GO TO 150	NEW	1220
C	LOOP THROUGH THE INTERPOLATATION POINTS	NEW	1230
	DO 140 J=1,NI	NEW	1240
	IDIF = 2	NEW	1250
	NRLJ = NRL + J	NEW	1260
	DDTEMP(1,NRLJ) = FINTRP(J)	NEW	1270
C	LOOP THROUGH THE DIFFERENCES (IDIF INDEX)	NEW	1280
	NRLJ1 = NRLJ - 1	NEW	1290
	DO 130 K=1,NRLJ1	NEW	1300
	IPT = NRLJ - K	NEW	1310
	DIFFF = DDTEMP(IDIF-1,IPT+1) - DDTEMP(IDIF-1,IPT)	NEW	1320
	DIFFX = XDD(NRLJ) - XDD(IPT)	NEW	1330
	DDTEMP(IDIF,IPT) = DIFFF/DIFFX	NEW	1340
	IDIF = IDIF + 1	NEW	1350
130	CONTINUE	NEW	1360
140	CONTINUE	NEW	1370
150	CONTINUE	NEW	1380
	RETURN	NEW	1390
99999	FORMAT (9X, 30HNR2,NL2,IDIF,IPT,DIFFF,DIFFX =, 4I3, 2F12.6)	NEW	1400
	END	NEW	1410
	REAL FUNCTION POLYDL(X)	POL	10
C		POL	20
C	=====	POL	30
C		POL	40
C **	THIS FUNCTION EVALUATES THE CURRENT POLYNOMIAL PIECE REPRESENTED	POL	50

C	BY THE DIVIDED DIFFERENCES DDTEMP ON THE POINT SET XDD.	POL	60
C		POL	70
	REAL FUNCTION PPOLY(T, XKNOTS, COEFS, KDIMEN, NDIMEN)	PPO	10
C		PPO	20
C	-----	PPO	30
C		PPO	40
C	** THIS FUNCTION EVALUATES THE PIECEWISE POLYNOMIAL APPROXIMATION	PPO	50
C	COMPUTED IN ADAPT	PPO	60
C		PPO	70
	SUBROUTINE PTRANS(D, POWERS)	PTR	10
C		PTR	20
C	-----	PTR	30
C		PTR	40
C	** THIS PROGRAM CONVERTS POLYNOMIAL REPRESENTATION FROM DIVIDED	PTR	50
C	DIFFERENCE TO POWER FORM. THERE ARE COALESCED POINTS ON EACH	PTR	60
C	END OF THE INTERVAL (XL,XR) = (XLEFT(NSTACK),XRIGHT(NSTACK)).	PTR	70
C	THE NUMBER COALESCED AT EACH END IS LEFTX AND RIGHTX.	PTR	80
C	AND THERE ARE NINTRP OTHER PTS XINTRP(K) INBETWEEN THEM.	PTR	90
C	SEE SUBROUTINE NEWTON FOR MORE DETAILS	PTR	100
C		PTR	110
	STARTING REPRESENTATION IS (ASSUMING XL = 0)	PTR	450
C		PTR	460
C	D(1) +D(2)X +D(3)X**2 + --- +D(NL)X**(NL-1)	PTR	470
C	+(X**NL)*(D(NL+1) (+D(NL+2) (X-XR)**2 + --- +D(NL+NR)*(X-XR)**(NR-1)	PTR	480
C	*((X-XR)**NR)*(D(NL+NR+1) + D(NL+NR+2)*(X-XINTRP(1))	PTR	490
C	+D(NL+NR+3)*(X-XINTRP(1))(X-XINTRP(2)) + ---))	PTR	500
C		PTR	510
C	STRATEGY IS TO FIRST CONVERT THE PART FROM THE INTERP. PTS.	PTR	520
C	TO POLY IN (X-XR). THIS POLY THEN HAS ORIGIN SHIFTED TO XL.	PTR	530
C		PTR	540
C	THE CONVERSION OF THE INTERP PART IS DONE EXPLICITLY FOR DEGREE	PTR	550
C	TWO OR LESS AND DONE BY SYNTHETIC DIVISION FOR HIGHER DEGREES	PTR	560
C		PTR	570
C	D1 + D2(X-X1) +D3(X**2-(X1+X2)X +X1*X2)	PTR	580
C		PTR	590
C	THE RESULTING COEFFICIENTS ARE PUT IN THE ARRAY POWERS	PTR	600
C		PTR	610
	SUBROUTINE PUT(INTERV, XKNOTS, COEFS, KDIMEN, NDIMEN)	PUT	10
C		PUT	20
C	-----	PUT	30
C	** THIS PROGRAM PUTS INTERVALS ON THE STACK OR DISCARDS THEM.	PUT	40
C	WHEN AN INTERVAL IS DISCARDED A NEW KNOT IS FOUND. THEN THIS	PUT	50
C	PROGRAM UPDATES THE ERROR ESTIMATE, THE XKNOT ARRAY, TRANSFORMS	PUT	60
C	THE POLYNOMIAL TO THE POWER FORM AND PUT THE COEFFICIENTS INTO	PUT	70
C	THE ARRAY COEFS. IT ALSO CHECKS FOR PASSING BREAK POINTS	PUT	80
C		PUT	90
	SUBROUTINE SETUP(XKNOTS, COEFS, KDIMEN, NDIMEN)	SET	10
C		SET	20
C	-----	SET	30
C		SET	40
C	** THIS PROGRAM CHECKS THE INPUT DATA AND INITIALIZES THE COMPUTATION	SET	50
C		SET	60
	SUBROUTINE SUMARY(XKNOTS, COEFS, KDIMEN, NDIMEN)	SUM	10
C		SUM	20
C	-----	SUM	30
C		SUM	40
C	** THIS PROGRAM PRINTS OUT A SUMMARY OF RESULTS OF ADAPT	SUM	50
C		SUM	60
	SUBROUTINE TAKE(INTERV)	TAK	10
C		TAK	20
C	-----	TAK	30
C		TAK	40
C	** THIS PROGRAM TAKES AN ACTIVE INTERVAL OFF THE TOP OF THE STACK	TAK	50

C	IT ALSO DOES MOST OF THE WORK OF LOCATING AND HANDLING	TAK	60
C	BREAK POINTS	TAK	70
C		TAK	80
C	SUBROUTINE TERMIN(TEST, AND, PRINT, XKNOTS, KDIMEN)	TER	10
C		TER	20
C	-----	TER	30
C		TER	40
C	** THIS PROGRAM TESTS FOR TERMINATION AND GIVES INTERMEDIATE OUTPUT	TER	50
C		TER	60
	BLOCK DATA	BLK	10
C	SET PARAMETERS FOR FUNCTIONS IN F	BLK	20
	REAL PAR2, PAR3A, PAR3B, PAR4, PAR4A, P5A, P5B, P6A, P6B, P7,	BLK	30
	* P8A, P8B, P8C, P8D, P8E, P8F, P8G, P8H, P9, P10A, P10B, P10C	BLK	40
	COMMON /FPARS/ PAR2, PAR3A, PAR3B, PAR4, PAR4A, P5A, P5B, P6A,	BLK	50
	* P6B, P7, P8A, P8B, P8C, P8D, P8E, P8F, P8G, P8H, P9, P10A, P10B,	BLK	60
	* P10C	BLK	70
	DATA PAR2, PAR3A, PAR3B, PAR4, PAR4A, P5A, P5B	BLK	80
	* /.02,.4,.6,.001,6.,8000.,2037./	BLK	90
	DATA P6A, P6B, P7, P8A, P8B, P8C, P8D, P8E /4000.,3998.,.0001,-2.,	BLK	100
	* 1.5,7.,5.,9.6/	BLK	110
	DATA P8F, P8G, P8H, P9, P10A, P10B, P10C /.6,14.,16.,1.0.,25.,5,	BLK	120
	* .6667/	BLK	130
	END	BLK	140
C		MAN	10
C	**** THIS PROGRAM RUNS ADAPT WITH DATA READ IN FROM CARDS	MAN	20
C	IT IS DESIGNED TO TEST ADAPT WITH THE 20 FUNCTIONS IN F(X).	MAN	30
C	IT READS DATA FOR EVERYTHING EXCEPT EDIST(=ATYPE) AND PRINTS	MAN	40
C	A HEADING WITH THE FUNCTIONS NAME. IT CALLS AND TIMES ADAPT,	MAN	50
C	THEN INDEPENDENTLY CHECKS THE ERROR IN ADAPT AND PLOTS THE	MAN	60
C	FUNCTION AND ERROR CURVE.	MAN	70
C		MAN	80
C	NOTES -- USER MUST SUPPLY SYSTEM ROUTINES SECOND AND GRAPH	MAN	90
C	INTEGRATION USES ERRINT FROM ADAPT.	MAN	100
C	PRESENT DIMENSIONS (ON ZKNOTS + ZCOEFS) PREVENT USING	MAN	110
C	ERRINT TO CHECK ACCURACY FOR MORE THAN 100 KNOTS OR	MAN	120
C	POLYNOMIAL DEGREE 6.	MAN	130
C	QPOLY USED BY ERRINT IS A SINGLE ARGUMENT VERSION OF PPOLY.	MAN	140
C	SEE COMMENTS IN ADAPT ABOUT DIMENSION CONSTRAINTS AND	MAN	150
C	PORTABILITY.	MAN	160
C	IF ADAPT AND THIS DRIVER ARE CHANGED TO DOUBLE PRECISION	MAN	170
C	ONE PROBABLY WANTS TO LEAVE TIME,TSTART,TSTOP,XVALU,	MAN	180
C	GRAF1,GRAF2 AS SINGLE PRECISION	MAN	190
C	THEY ARE SEPARATELY DECLARED HERE.	MAN	200
C		MAN	210
	REAL FUNCTION F(X, FDERV)	F	10
C		F	20
C	** AN ARRAY OF TWENTY TEST FUNCTIONS WITH 2 TO 5 DERIVATIVES.	F	30
C	THE DERIVATIVE VALUES ARE PLACED IN FDERV - DIMENSIONED AT 5	F	40
C	THE FUNCTIONS OF THE SECOND GROUP ARE PARAMETERIZED IN VARIOUS	F	50
C	WAYS. THE PARAMETERS ARE SET IN BLOCK DATA AND AVAILABLE	F	60
C	THROUGH THE COMMON BLOCK / FPARS /	F	70
C	REQUIRED INPUT INFORMATION IN COMMON BLOCK / FDATA /	F	80
C	JFUNK = INDEX OF THE FUNCTION SELECTED	F	90
C	KOUNT = NUMBER OF F-EVALUATIONS	F	100
C	MAXK = MAXIMUM ALLOWED VALUE OF KOUNT	F	110
C	REQUIRED CONTROL INFORMATION IN COMMON BLOCK / KONTROL /	F	120
C	FINISH = SWITCH THAT STOPS ADAPT	F	130
C		F	140

ALGORITHM 526

Bivariate Interpolation and Smooth Surface Fitting for Irregularly Distributed Data Points [E1]

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Key Words and Phrases: bivariate interpolation, interpolation, partial derivative, polynomial, smooth surface fitting

CR Categories: 5.13

Language: ANSI Standard Fortran

DESCRIPTION

This algorithm describes the IDBVIP/IDSFFT subprogram package that implements the method of bivariate interpolation and smooth surface fitting for irregularly distributed data points [1]. The package is written in ANSI Standard Fortran [2].

The package consists of nine subprograms, i.e. eight subroutines and a function. Two subroutines, IDBVIP and IDSFFT, are the master subroutines of the package, and each interfaces with the user. The IDBVIP subroutine performs bivariate interpolation; it estimates the z values at the specified points in the x - y plane. The IDSFFT subroutine performs smooth surface fitting; it estimates the z values at the specified rectangular grid points in the x - y plane and generates a doubly-dimensioned array containing these estimated values.

The remaining six subroutines are supporting subroutines called by either IDBVIP or IDSFFT or by both: IDCLDP determines several data points closest to each of the data points; IDGRID organizes output grid points for IDSFFT by sorting them in ascending order of triangle numbers; IDLCTN locates a point (i.e. determines a triangle in which the point lies) for IDBVIP; IDPDRV estimates partial derivatives at the data points; IDPTIP performs punctual interpolation (i.e. interpolation at a point); and IDTANG triangulates (or divides into a number of triangles) the x - y plane. A function, IDXCHG, is called by IDTANG and determines whether or not an exchange of triangles is necessary.

The package includes two common blocks, IDLC and IDPI. Including these common areas, the package occupies approximately 3500 locations on the CDC-6600 computer.

When the user wishes to call either the IDBVIP or the IDSFFT subroutine repeatedly with identical data as parts of input data in two consecutive calls, he can save computation times considerably by specifying an appropriate mode of computation. (This mode is specified with the MD parameter in the call statements.)

Tables I and II show the approximate computation times required for the IDBVIP and IDSFFT subroutines, respectively, on the CDC-6600 computer when

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TABLE I. Approximate Computation Times Required for the IDBVIP Subroutine on the CDC-6600 Computer

NDP	NIP	Time (seconds)		
		MD = 1	MD = 2	MD = 3
10	10	.03	.02	.02
	100	.13	.12	.08
	1000	1.1	1.1	.70
100	10	.75	.10	.07
	100	.86	.22	.14
	1000	2.1	1.5	.84
1000	10	47.	.84	.60
	100	47.	1.1	.68
	1000	50.	3.8	1.4

TABLE II. Approximate Computation Times Required for the IDSFFT Subroutine on the CDC-6600 Computer

NDP	NXI * NYI	Time (seconds)		
		MD = 1	MD = 2	MD = 3
10	11 * 11	.08	.07	.04
	33 * 33	.41	.40	.22
	101 * 101	3.5	3.4	1.8
100	11 * 11	.86	.20	.14
	33 * 33	1.3	.66	.37
	101 * 101	4.6	4.0	2.0
1000	11 * 11	47.	1.1	.70
	33 * 33	49.	2.6	1.3
	101 * 101	54.	8.2	3.5

four additional data points are used for estimating partial derivatives at each data point (i.e. NCP = 4). In these tables, NDP is the number of data points, NIP is the number of output points at which interpolation is to be performed, and NXI and NYI are the numbers of output grid points in the x and y coordinates. MD is the mode of computation; MD = 1 for a new NCP value and/or a new set of x - y coordinates of data points, MD = 2 for an old NCP value, an old set of x - y coordinates of data points, and a new set of x - y coordinates of output points, and MD = 3 for an old NCP value, an old set of x - y coordinates of data points, and an old set of x - y coordinates of output points.

REFERENCES

1. AKIMA, H. A method of bivariate interpolation and smooth surface fitting for irregularly distributed data points. *ACM Trans. Math. Software* 4, 2 (June 1978), 148-159.
2. ANSI Standard Fortran. Pub. X3.9-1966. Amer. Nat. Standards Inst., New York; also reproduced in W.P. Heising's history and summary of FORTRAN standardization development for the ASA. *Comm. ACM* 7, 10 (Oct. 1964), 590-625.

ALGORITHM

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C      PROGRAM TTIDBS(OUTPUT,TAPE6=OUTPUT)
C THIS PROGRAM IS A TEST PROGRAM FOR THE IDBVIP/IDSFFT SUBPRO-
C GRAM PACKAGE. ALL ELEMENTS OF RESULTING DZI1 AND DZI2 ARRAYS
C ARE EXPECTED TO BE ZERO.
C THE LUN CONSTANT IN THE LAST DATA INITIALIZATION STATEMENT IS
C THE LOGICAL UNIT NUMBER OF THE STANDARD OUTPUT UNIT AND IS,
C THEREFORE, SYSTEM DEPENDENT.
C DECLARATION STATEMENTS
      DIMENSION XD(30),YD(30),ZD(30),
1             XI(6),YI(5),ZI(6,5),
2             ZI1(6,5),ZI2(6,5),DZI1(6,5),DZI2(6,5),
3             IWK(1030),WK(240)
      DATA NCP/4/
      DATA NDP/30/
      DATA XD(1),XD(2),XD(3),XD(4),XD(5),XD(6),
1         XD(7),XD(8),XD(9),XD(10),XD(11),XD(12),
2         XD(13),XD(14),XD(15),XD(16),XD(17),XD(18),
3         XD(19),XD(20),XD(21),XD(22),XD(23),XD(24),
4         XD(25),XD(26),XD(27),XD(28),XD(29),XD(30)/
5         11.16,24.20,19.85,10.35,19.72,0.00,
6         20.87,19.99,10.28,4.51,0.00,16.70,
7         6.08,25.00,14.90,0.00,9.66,5.22,
8         11.77,15.10,25.00,25.00,14.59,15.20,
9         5.23,2.14,0.51,25.00,21.67,3.31/
      DATA YD(1),YD(2),YD(3),YD(4),YD(5),YD(6),
1         YD(7),YD(8),YD(9),YD(10),YD(11),YD(12),
2         YD(13),YD(14),YD(15),YD(16),YD(17),YD(18),
3         YD(19),YD(20),YD(21),YD(22),YD(23),YD(24),
4         YD(25),YD(26),YD(27),YD(28),YD(29),YD(30)/
5         1.24,16.23,10.72,4.11,1.39,20.00,
6         20.00,4.62,15.16,20.00,4.48,19.65,
7         4.58,11.87,3.12,0.00,20.00,14.66,
8         10.47,17.19,3.87,0.00,8.71,0.00,
9         10.72,15.03,8.37,20.00,14.36,0.13/
      DATA ZD(1),ZD(2),ZD(3),ZD(4),ZD(5),ZD(6),
1         ZD(7),ZD(8),ZD(9),ZD(10),ZD(11),ZD(12),
2         ZD(13),ZD(14),ZD(15),ZD(16),ZD(17),ZD(18),
3         ZD(19),ZD(20),ZD(21),ZD(22),ZD(23),ZD(24),
4         ZD(25),ZD(26),ZD(27),ZD(28),ZD(29),ZD(30)/
5         22.15,2.83,7.97,22.33,16.83,34.60,
6         5.74,14.72,21.59,15.61,61.77,6.31,
7         35.74,4.40,21.70,58.20,4.73,40.36,
8         13.62,12.57,8.74,12.00,14.81,21.60,
9         26.50,53.10,49.43,0.60,5.52,44.08/
      DATA NXI/6/,NYI/5/
      DATA XI(1),XI(2),XI(3),XI(4),XI(5),XI(6)/
1         0.00,5.00,10.00,15.00,20.00,25.00/
      DATA YI(1),YI(2),YI(3),YI(4),YI(5)/
1         0.00,5.00,10.00,15.00,20.00/
      DATA ZI(1,1),ZI(2,1),ZI(3,1),ZI(4,1),ZI(5,1),ZI(6,1),
1         ZI(1,2),ZI(2,2),ZI(3,2),ZI(4,2),ZI(5,2),ZI(6,2),
2         ZI(1,3),ZI(2,3),ZI(3,3),ZI(4,3),ZI(5,3),ZI(6,3),
3         ZI(1,4),ZI(2,4),ZI(3,4),ZI(4,4),ZI(5,4),ZI(6,4),
4         ZI(1,5),ZI(2,5),ZI(3,5),ZI(4,5),ZI(5,5),ZI(6,5)/
5         58.20,39.55,26.90,21.71,17.68,12.00,
6         61.58,39.39,22.04,21.29,14.36,8.04,
7         59.18,27.39,16.78,13.25,8.59,5.36,
8         52.82,40.27,22.76,16.61,7.40,2.88,
9         34.60,14.05,4.12,3.17,6.31,0.60/
      DATA LUN/6/
C CALCULATION
10 MD=1
      DO 12 IYI=1,NYI
          DO 11 IXI=1,NXI
              IF(IXI.NE.1.OR.IYI.NE.1) MD=2
              CALL IDBVIP(MD,NCP,NDP,XD,YD,ZD,1,XI(IXI),YI(IYI),
1                 ZI1(IXI,IYI),IWK,WK)
11          CONTINUE
12      CONTINUE
15 CALL IDSFFT(1,NCP,NDP,XD,YD,ZD,NXI,NYI,XI,YI,ZI2,IWK,WK)
      DO 17 IYI=1,NYI
          DO 16 IXI=1,NXI
              DZI1(IXI,IYI)=ABS(ZI1(IXI,IYI)-ZI(IXI,IYI))
              DZI2(IXI,IYI)=ABS(ZI2(IXI,IYI)-ZI(IXI,IYI))

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16 CONTINUE
17 CONTINUE
C PRINTING OF INPUT DATA
20 WRITE (LUN,2020) NDP
DO 23 IDP=1,NDP
  IF (MOD(IDP,5).EQ.1) WRITE (LUN,2021)
  WRITE (LUN,2022) IDP,XD(IDP),YD(IDP),ZD(IDP)
23 CONTINUE
C PRINTING OF OUTPUT RESULTS
30 WRITE (LUN,2030)
WRITE (LUN,2031) YI
DO 33 IXI=1,NXI
  WRITE (LUN,2032) XI(IXI),(Z11(IXI,IYI),IYI=1,NYI)
33 CONTINUE
40 WRITE (LUN,2040)
WRITE (LUN,2031) YI
DO 43 IXI=1,NXI
  WRITE (LUN,2032) XI(IXI),(DZ11(IXI,IYI),IYI=1,NYI)
43 CONTINUE
50 WRITE (LUN,2050)
WRITE (LUN,2031) YI
DO 53 IXI=1,NXI
  WRITE (LUN,2032) XI(IXI),(Z12(IXI,IYI),IYI=1,NYI)
53 CONTINUE
60 WRITE (LUN,2060)
WRITE (LUN,2031) YI
DO 63 IXI=1,NXI
  WRITE (LUN,2032) XI(IXI),(DZ12(IXI,IYI),IYI=1,NYI)
63 CONTINUE
STOP
C FORMAT STATEMENTS
2020 FORMAT(1H1,6HTTIDBS/////3X,10HINPUT DATA,8X,5HNDP =,I3//
1 30H I XD YD ZD /)
2021 FORMAT(1X)
2022 FORMAT(5X,I2,2X,3F7.2)
2030 FORMAT(1H1,6HTTIDBS/////3X,17HIDBVIP SUBROUTINE//
1 26X,10HZ11(XI,YI))
2031 FORMAT(7X,2HXI,4X,3HYI=/12X,5F7.2/)
2032 FORMAT(1X/1X,F9.2,2X,5F7.2)
2040 FORMAT(1X/////3X,10HDIFFERENCE//
1 25X,11HDZ11(XI,YI))
2050 FORMAT(1H1,6HTTIDBS/////3X,17HIDSFFT SUBROUTINE//
1 26X,10HZ12(XI,YI))
2060 FORMAT(1X/////3X,10HDIFFERENCE//
1 25X,11HDZ12(XI,YI))
END

SUBROUTINE IDBVIP(MD,NCP,NDP,XD,YD,ZD,NIP,XI,YI,ZI,
1 IWK,WK)
C THIS SUBROUTINE PERFORMS BIVARIATE INTERPOLATION WHEN THE PRO-
C JECTIONS OF THE DATA POINTS IN THE X-Y PLANE ARE IRREGULARLY
C DISTRIBUTED IN THE PLANE.
C THE INPUT PARAMETERS ARE
C MD = MODE OF COMPUTATION (MUST BE 1, 2, OR 3),
C = 1 FOR NEW NCP AND/OR NEW XD-YD,
C = 2 FOR OLD NCP, OLD XD-YD, NEW XI-YI,
C = 3 FOR OLD NCP, OLD XD-YD, OLD XI-YI,
C NCP = NUMBER OF ADDITIONAL DATA POINTS USED FOR ESTI-
C MATING PARTIAL DERIVATIVES AT EACH DATA POINT
C (MUST BE 2 OR GREATER, BUT SMALLER THAN NDP),
C NDP = NUMBER OF DATA POINTS (MUST BE 4 OR GREATER),
C XD = ARRAY OF DIMENSION NDP CONTAINING THE X
C COORDINATES OF THE DATA POINTS,
C YD = ARRAY OF DIMENSION NDP CONTAINING THE Y
C COORDINATES OF THE DATA POINTS,
C ZD = ARRAY OF DIMENSION NDP CONTAINING THE Z
C COORDINATES OF THE DATA POINTS,
C NIP = NUMBER OF OUTPUT POINTS AT WHICH INTERPOLATION
C IS TO BE PERFORMED (MUST BE 1 OR GREATER),
C XI = ARRAY OF DIMENSION NIP CONTAINING THE X
C COORDINATES OF THE OUTPUT POINTS,
C YI = ARRAY OF DIMENSION NIP CONTAINING THE Y
C COORDINATES OF THE OUTPUT POINTS.
C THE OUTPUT PARAMETER IS

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```

ID000810
ID000820
ID000830
ID000840
ID000850
ID000860
ID000870
ID000880
ID000890
ID000900
ID000910
ID000920
ID000930
ID000940
ID000950
ID000960
ID000970
ID000980
ID000990
ID001000
ID001010
ID001020
ID001030
ID001040
ID001050
ID001060
ID001070
ID001080
ID001090
ID001100
ID001110
ID001120
ID001130
ID001140
ID001150
ID001160
ID001170
ID001180
ID001190
ID001200
ID001210
ID001220
ID001230
ID001240
ID001250
ID001260

ID001340
ID001350
ID001360
ID001370
ID001380
ID001390
ID001400
ID001410
ID001420
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ID001470
ID001480
ID001490
ID001500
ID001510
ID001520
ID001530
ID001540
ID001550
ID001560
ID001570
ID001580
ID001590
ID001600

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C      ZI = ARRAY OF DIMENSION NIP WHERE INTERPOLATED Z          ID001610
C      VALUES ARE TO BE STORED.                                ID001620
C THE OTHER PARAMETERS ARE                                     ID001630
C      IWK = INTEGER ARRAY OF DIMENSION                        ID001640
C      MAX0(31,27+NCP)*NDP+NIP                                ID001650
C      USED INTERNALLY AS A WORK AREA,                          ID001660
C      WK = ARRAY OF DIMENSION 8*NDP USED INTERNALLY AS A     ID001670
C      WORK AREA.                                              ID001680
C THE VERY FIRST CALL TO THIS SUBROUTINE AND THE CALL WITH A NEW
C NCP VALUE, A NEW NDP VALUE, AND/OR NEW CONTENTS OF THE XD AND
C YD ARRAYS MUST BE MADE WITH MD=1. THE CALL WITH MD=2 MUST BE
C PRECEDED BY ANOTHER CALL WITH THE SAME NCP AND NDP VALUES AND
C WITH THE SAME CONTENTS OF THE XD AND YD ARRAYS. THE CALL WITH
C MD=3 MUST BE PRECEDED BY ANOTHER CALL WITH THE SAME NCP, NDP,
C AND NIP VALUES AND WITH THE SAME CONTENTS OF THE XD, YD, XI,
C AND YI ARRAYS. BETWEEN THE CALL WITH MD=2 OR MD=3 AND ITS
C PRECEDING CALL, THE IWK AND WK ARRAYS MUST NOT BE DISTURBED.
C USE OF A VALUE BETWEEN 3 AND 5 (INCLUSIVE) FOR NCP IS RECOM-
C MENDED UNLESS THERE ARE EVIDENCES THAT DICTATE OTHERWISE.   ID001790
C THE LUN CONSTANT IN THE DATA INITIALIZATION STATEMENT IS THE
C LOGICAL UNIT NUMBER OF THE STANDARD OUTPUT UNIT AND IS,      ID001800
C THEREFORE, SYSTEM DEPENDENT.                                ID001810
C THIS SUBROUTINE CALLS THE IDCLDP, IDLCTN, IDPDRV, IDPTIP, AND
C IDTANG SUBROUTINES.                                         ID001820
C DECLARATION STATEMENTS                                     ID001830
C      DIMENSION XD(100),YD(100),ZD(100),XI(1000),YI(1000),
C      1      ZI(1000),IWK(4100),WK(800)
C      COMMON/IDLC/NIT
C      COMMON/IDPI/ITPV
C      DATA LUN/6/
C SETTING OF SOME INPUT PARAMETERS TO LOCAL VARIABLES.
C (FOR MD=1,2,3)
C      MD0=MD
C      NCP0=NCP
C      NDP0=NDP
C      NIP0=NIP
C ERROR CHECK. (FOR MD=1,2,3)
C      20 IF (MD0.LT.1.OR.MD0.GT.3) GO TO 90
C      IF (NCP0.LT.2.OR.NCP0.GE.NDP0) GO TO 90
C      IF (NDP0.LT.4) GO TO 90
C      IF (NIP0.LT.1) GO TO 90
C      IF (MD0.GE.2) GO TO 21
C      IWK(1)=NCP0
C      IWK(2)=NDP0
C      GO TO 22
C      21 NCPPV=IWK(1)
C      NDPPV=IWK(2)
C      IF (NCP0.NE.NCPPV) GO TO 90
C      IF (NDP0.NE.NDPPV) GO TO 90
C      22 IF (MD0.GE.3) GO TO 23
C      IWK(3)=NIP
C      GO TO 30
C      23 NIPPV=IWK(3)
C      IF (NIP0.NE.NIPPV) GO TO 90
C ALLOCATION OF STORAGE AREAS IN THE IWK ARRAY. (FOR MD=1,2,3)
C      30 JWIPT=16
C      JWIWL=6*NDP0+1
C      JWIWK=JWIWL
C      JWIPL=24*NDP0+1
C      JWIWP=30*NDP0+1
C      JWIPC=27*NDP0+1
C      JWIT0=MAX0(31,27+NCP0)*NDP0
C TRIANGULATES THE X-Y PLANE. (FOR MD=1)
C      40 IF (MD0.GT.1) GO TO 50
C      CALL IDTANG (NDP0,XD,YD,NT,IWK(JWIPT),NL,IWK(JWIPL),
C      1      IWK(JWIWL),IWK(JWIWP),WK)
C      IWK(5)=NT
C      IWK(6)=NL
C      IF (NT.EQ.0) RETURN
C DETERMINES NCP POINTS CLOSEST TO EACH DATA POINT. (FOR MD=1)
C      50 IF (MD0.GT.1) GO TO 60
C      CALL IDCLDP (NDP0,XD,YD,NCP0,IWK(JWIPC))
C      IF (IWK(JWIPC).EQ.0) RETURN
C LOCATES ALL POINTS AT WHICH INTERPOLATION IS TO BE PERFORMED.
C (FOR MD=1,2)
C      60 IF (MD0.EQ.3) GO TO 70
C      NIT=0

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        JWIT=JWIT0
        DO 61 IIP=1,NIP0
            JWIT=JWIT+1
            CALL IDLCTN(NDP0,XD,YD,NT,IWK(.TWIPT),NL,IWK(JWIPL),
1             XI(IIP),YI(IIP),IWK(JWIT),IWK(JWIWK),WK)
        61 CONTINUE
C ESTIMATES PARTIAL DERIVATIVES AT ALL DATA POINTS.
C (FOR MD=1,2,3)
        70 CALL IDPDRV(NDP0,XD,YD,ZD,NCP0,IWK(JWIPC),WK)
C INTERPOLATES THE ZI VALUES. (FOR MD=1,2,3)
        80 ITPV=0
            JWIT=JWIT0
            DO 81 IIP=1,NIP0
                JWIT=JWIT+1
                CALL IDPTIP(XD,YD,ZD,NT,IWK(JWIPT),NL,IWK(JWIPL),WK,
1             IWK(JWIT),XI(IIP),YI(IIP),ZI(IIP))
            81 CONTINUE
        RETURN
C ERROR EXIT
        90 WRITE (LUN,2090) MD0,NCP0,NDP0,NIP0
        RETURN
C FORMAT STATEMENT FOR ERROR MESSAGE
        2090 FORMAT(1X/4H *** IMPROPER INPUT PARAMETER VALUE(S) /
1     7H MD =,I4,10X,5HNCP =,I6,10X,5HNDP =,I6,
2     10X,5HNIP =,I6/
3     35H ERROR DETECTED IN ROUTINE IDBVIP/)
        END

        SUBROUTINE IDCLDP(NDP,XD,YD,NCP,IPC)
C THIS SUBROUTINE SELECTS SEVERAL DATA POINTS THAT ARE CLOSEST
C TO EACH OF THE DATA POINT.
C THE INPUT PARAMETERS ARE
C   NDP = NUMBER OF DATA POINTS,
C   XD,YD = ARRAYS OF DIMENSION NDP CONTAINING THE X AND Y
C           COORDINATES OF THE DATA POINTS,
C   NCP = NUMBER OF DATA POINTS CLOSEST TO EACH DATA
C           POINTS.
C THE OUTPUT PARAMETER IS
C   IPC = INTEGER ARRAY OF DIMENSION NCP*NDP, WHERE THE
C         POINT NUMBERS OF NCP DATA POINTS CLOSEST TO
C         EACH OF THE NDP DATA POINTS ARE TO BE STORED.
C THIS SUBROUTINE ARBITRARILY SETS A RESTRICTION THAT NCP MUST
C NOT EXCEED 25.
C THE LUN CONSTANT IN THE DATA INITIALIZATION STATEMENT IS THE
C LOGICAL UNIT NUMBER OF THE STANDARD OUTPUT UNIT AND IS,
C THEREFORE, SYSTEM DEPENDENT.
C DECLARATION STATEMENTS
        DIMENSION XD(100),YD(100),IPC(400)
        DIMENSION DSQ0(25),JPC0(25)
        DATA NCPMX/25/,LUN/6/
C STATEMENT FUNCTION
        DSQF(U1,V1,U2,V2)=(U2-U1)**2+(V2-V1)**2
C PRELIMINARY PROCESSING
        10 NDP0=NDP
            NCP0=NCP
            IF(NDP0.LT.2) GO TO 90
            IF(NCP0.LT.1.OR.NCP0.GT.NCPMX.OR.NCP0.GE.NDP0) GO TO 90
C CALCULATION
        20 DO 59 IP1=1,NDP0
C - SELECTS NCP POINTS.
            X1=XD(IP1)
            Y1=YD(IP1)
            J1=0
            DSQMX=0.0
            DO 22 IP2=1,NDP0
                IF(IP2.EQ.IP1) GO TO 22
                DSQI=DSQF(X1,Y1,XD(IP2),YD(IP2))
                J1=J1+1
                DSQ0(J1)=DSQI
                IPC0(J1)=IP2
                IF(DSQI.LE.DSQMX) GO TO 21
                DSQMX=DSQI
                JMX=J1
            21 IF(J1.GE.NCP0) GO TO 23
            22 CONTINUE

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ID002380
ID002390
ID002400
ID002410
ID002420
ID002430
ID002440
ID002450
ID002460
ID002470
ID002480
ID002490
ID002500
ID002510
ID002520
ID002530
ID002540
ID002550
ID002560
ID002570
ID002580
ID002590
ID002600
ID002610
ID002620
ID002630
ID002640

ID002720
ID002730
ID002740
ID002750
ID002760
ID002770
ID002780
ID002790
ID002800
ID002810
ID002820
ID002830
ID002840
ID002850
ID002860
ID002870
ID002880
ID002890
ID002900
ID002910
ID002920
ID002930
ID002940
ID002950
ID002960
ID002970
ID002980
ID002990
ID003000
ID003010
ID003020
ID003030
ID003040
ID003050
ID003060
ID003070
ID003080
ID003090
ID003100
ID003110
ID003120
ID003130
ID003140
ID003150
ID003160
ID003170
ID003180

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23  IP2MN=IP2+1                                ID003190
    IF(IP2MN.GT.NDP0)      GO TO 30            ID003200
    DO 25  IP2=IP2MN,NDP0  ID003210
        IF(IP2.EQ.IP1)    GO TO 25            ID003220
        DSQI=DSQF(X1,Y1,XD(IP2),YD(IP2))     ID003230
        IF(DSQI.GE.DSQMX) GO TO 25            ID003240
        DSQ0(JMX)=DSQI    ID003250
        IPC0(JMX)=IP2     ID003260
        DSQMX=0.0         ID003270
        DO 24  J1=1,NCP0  ID003280
            IF(DSQ0(J1).LE.DSQMX) GO TO 24    ID003290
            DSQMX=DSQ0(J1) ID003300
            JMX=J1        ID003310
24  CONTINUE                                ID003320
25  CONTINUE                                ID003330
C - CHECKS IF ALL THE NCP+1 POINTS ARE COLLINEAR. ID003340
30  IP2=IPC0(1)                               ID003350
    DX12=XD(IP2)-X1                            ID003360
    DY12=YD(IP2)-Y1                            ID003370
    DO 31  J3=2,NCP0                            ID003380
        IP3=IPC0(J3)                            ID003390
        DX13=XD(IP3)-X1                        ID003400
        DY13=YD(IP3)-Y1                        ID003410
        IF((DY13*DX12-DX13*DY12).NE.0.0)      GO TO 50 ID003420
31  CONTINUE                                ID003430
C - SEARCHES FOR THE CLOSEST NONCOLLINEAR POINT. ID003440
40  NCLPT=0                                    ID003450
    DO 43  IP3=1,NDP0                            ID003460
        IF(IP3.EQ.IP1)      GO TO 43            ID003470
        DO 41  J4=1,NCP0  ID003480
            IF(IP3.EQ.IPC0(J4)) GO TO 43        ID003490
41  CONTINUE                                ID003500
    DX13=XD(IP3)-X1                            ID003510
    DY13=YD(IP3)-Y1                            ID003520
    IF((DY13*DX12-DX13*DY12).EQ.0.0)          GO TO 43 ID003530
    DSQI=DSQF(X1,Y1,XD(IP3),YD(IP3))           ID003540
    IF(NCLPT.EQ.0)      GO TO 42                ID003550
    IF(DSQI.GE.DSQMN)   GO TO 43                ID003560
42  NCLPT=1                                    ID003570
    DSQMN=DSQI    ID003580
    IP3MN=IP3    ID003590
43  CONTINUE                                ID003600
    IF(NCLPT.EQ.0)      GO TO 91                ID003610
    DSQMX=DSQMN    ID003620
    IPC0(JMX)=IP3MN ID003630
C - REPLACES THE LOCAL ARRAY FOR THE OUTPUT ARRAY. ID003640
50  J1=(IP1-1)*NCP0  ID003650
    DO 51  J2=1,NCP0  ID003660
        J1=J1+1  ID003670
        IPC(J1)=IPC0(J2) ID003680
51  CONTINUE                                ID003690
59  CONTINUE                                ID003700
    RETURN  ID003710
C ERROR EXIT                                ID003720
90  WRITE (LUN,2090)  ID003730
    GO TO 92  ID003740
91  WRITE (LUN,2091)  ID003750
92  WRITE (LUN,2092)  NDP0,NCP0  ID003760
    IPC(1)=0  ID003770
    RETURN  ID003780
C FORMAT STATEMENTS FOR ERROR MESSAGES      ID003790
2090 FORMAT(1X/41H ***  IMPROPER INPUT PARAMETER VALUE(S).) ID003800
2091 FORMAT(1X/33H ***  ALL COLLINEAR DATA POINTS.)  ID003810
2092 FORMAT(8H  NDP =,I5,5X,5HNCP =,I5/  ID003820
1 35H ERROR DETECTED IN ROUTINE  IDCLDF/)  ID003830
END  ID003840

SUBROUTINE IDGRID(XD, YD, NT, IPT, NL, IPL, NXI, NYI, XI, YI,
*  NGP, IGP)
C THIS SUBROUTINE ORGANIZES GRID POINTS FOR SURFACE FITTING BY
C SORTING THEM IN ASCENDING ORDER OF TRIANGLE NUMBERS AND OF THE
C BORDER LINE SEGMENT NUMBER.
C THE INPUT PARAMETERS ARE
C   XD,YD = ARRAYS OF DIMENSION NDP CONTAINING THE X AND Y
C   COORDINATES OF THE DATA POINTS, WHERE NDP IS THE

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IDG  10
IDG  20
IDG  30
IDG  40
IDG  50
IDG  60
IDG  70
IDG  80

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C          NUMBER OF THE DATA POINTS,          IDG  90
C  NT = NUMBER OF TRIANGLES,                    IDG 100
C  IPT = INTEGER ARRAY OF DIMENSION 3*NT CONTAINING THE IDG 110
C          POINT NUMBERS OF THE VERTEXES OF THE TRIANGLES, IDG 120
C  NL = NUMBER OF BORDER LINE SEGMENTS,         IDG 130
C  IPL = INTEGER ARRAY OF DIMENSION 3*NL CONTAINING THE IDG 140
C          POINT NUMBERS OF THE END POINTS OF THE BORDER IDG 150
C          LINE SEGMENTS AND THEIR RESPECTIVE TRIANGLE IDG 160
C          NUMBERS,                               IDG 170
C  NXI = NUMBER OF GRID POINTS IN THE X COORDINATE, IDG 180
C  NYI = NUMBER OF GRID POINTS IN THE Y COORDINATE, IDG 190
C  XI,YI = ARRAYS OF DIMENSION NXI AND NYI CONTAINING IDG 200
C          THE X AND Y COORDINATES OF THE GRID POINTS, IDG 210
C          RESPECTIVELY.                          IDG 220
C THE OUTPUT PARAMETERS ARE                      IDG 230
C  NGP = INTEGER ARRAY OF DIMENSION 2*(NT+2*NL) WHERE THE IDG 240
C          NUMBER OF GRID POINTS THAT BELONG TO EACH OF THE IDG 250
C          TRIANGLES OR OF THE BORDER LINE SEGMENTS ARE TO IDG 260
C          BE STORED,                             IDG 270
C  IGP = INTEGER ARRAY OF DIMENSION NXI*NYI WHERE THE IDG 280
C          GRID POINT NUMBERS ARE TO BE STORED IN ASCENDING IDG 290
C          ORDER OF THE TRIANGLE NUMBER AND THE BORDER LINE IDG 300
C          SEGMENT NUMBER.                        IDG 310
C DECLARATION STATEMENTS                       IDG 320
C   DIMENSION XD(100), YD(100), IPT(585), IPL(300), XI(101), IDG 330
C   *   YI(101), NGP(800), IGP(10201)           IDG 340
C STATEMENT FUNCTIONS                          IDG 350
C   SIDE(U1,V1,U2,V2,U3,V3) = (U1-U3)*(V2-V3) - (V1-V3)*(U2-U3) IDG 360
C   SPDT(U1,V1,U2,V2,U3,V3) = (U1-U2)*(U3-U2) + (V1-V2)*(V3-V2) IDG 370
C PRELIMINARY PROCESSING                      IDG 380
C   NT0 = NT                                    IDG 390
C   NL0 = NL                                    IDG 400
C   NXI0 = NXI                                  IDG 410
C   NYI0 = NYI                                  IDG 420
C   NXINYI = NXI0*NYI0                         IDG 430
C   XIMN = AMIN1(XI(1),XI(NXI0))                IDG 440
C   XIMX = AMAX1(XI(1),XI(NXI0))                IDG 450
C   YIMN = AMIN1(YI(1),YI(NYI0))                IDG 460
C   YIMX = AMAX1(YI(1),YI(NYI0))                IDG 470
C DETERMINES GRID POINTS INSIDE THE DATA AREA. IDG 480
C   JNGP0 = 0                                    IDG 490
C   JNGP1 = 2*(NT0+2*NL0) + 1                   IDG 500
C   JIGP0 = 0                                    IDG 510
C   JIGP1 = NXINYI + 1                          IDG 520
C   DO 160 IT0=1,NT0                             IDG 530
C     NGP0 = 0                                    IDG 540
C     NGP1 = 0                                    IDG 550
C     IT0T3 = IT0*3                              IDG 560
C     IP1 = IPT(IT0T3-2)                         IDG 570
C     IP2 = IPT(IT0T3-1)                         IDG 580
C     IP3 = IPT(IT0T3)                           IDG 590
C     X1 = XD(IP1)                               IDG 600
C     Y1 = YD(IP1)                               IDG 610
C     X2 = XD(IP2)                               IDG 620
C     Y2 = YD(IP2)                               IDG 630
C     X3 = XD(IP3)                               IDG 640
C     Y3 = YD(IP3)                               IDG 650
C     XMN = AMIN1(X1,X2,X3)                      IDG 660
C     XMX = AMAX1(X1,X2,X3)                      IDG 670
C     YMN = AMIN1(Y1,Y2,Y3)                      IDG 680
C     YMX = AMAX1(Y1,Y2,Y3)                      IDG 690
C     INSD = 0                                    IDG 700
C     DO 20 IXI=1,NXI0                            IDG 710
C       IF (XI(IXI).GE.XMN .AND. XI(IXI).LE.XMX) GO TO 10 IDG 720
C       IF (INSD.EQ.0) GO TO 20                   IDG 730
C       IXIMX = IXI - 1                           IDG 740
C       GO TO 30                                  IDG 750
10    IF (INSD.EQ.1) GO TO 20                     IDG 760
C     INSD = 1                                    IDG 770
C     IXIMN = IXI                                 IDG 780
20    CONTINUE                                    IDG 790
C     IF (INSD.EQ.0) GO TO 150                   IDG 800
C     IXIMX = NXI0                                IDG 810
30    DO 140 IYI=1,NYI0                          IDG 820
C     YII = YI(IYI)                              IDG 830
C     IF (YII.LT.YMN .OR. YII.GT.YMX) GO TO 140 IDG 840
C     DO 130 IXI=IXIMN,IXIMX                     IDG 850

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	XII = XI(IXI)	IDG 860
	L = 0	IDG 870
	IF (SIDE(X1,Y1,X2,Y2,XII,YII)) 130, 40, 50	IDG 880
40	L = 1	IDG 890
50	IF (SIDE(X2,Y2,X3,Y3,XII,YII)) 130, 60, 70	IDG 900
60	L = 1	IDG 910
70	IF (SIDE(X3,Y3,X1,Y1,XII,YII)) 130, 80, 90	IDG 920
80	L = 1	IDG 930
90	IZI = NXI0*(IYI-1) + IXI	IDG 940
	IF (L.EQ.1) GO TO 100	IDG 950
	NGP0 = NGP0 + 1	IDG 960
	JIGP0 = JIGP0 + 1	IDG 970
	IGP(JIGP0) = IZI	IDG 980
	GO TO 130	IDG 990
100	IF (JIGP1.GT.NXINYI) GO TO 120	IDG 1000
	DO 110 JIGP1I=JIGP1,NXINYI	IDG 1010
	IF (IZI.EQ.IGP(JIGP1I)) GO TO 130	IDG 1020
110	CONTINUE	IDG 1030
120	NGP1 = NGP1 + 1	IDG 1040
	JIGP1 = JIGP1 - 1	IDG 1050
	IGP(JIGP1) = IZI	IDG 1060
130	CONTINUE	IDG 1070
140	CONTINUE	IDG 1080
150	JNGP0 = JNGP0 + 1	IDG 1090
	NGP(JNGP0) = NGP0	IDG 1100
	JNGP1 = JNGP1 - 1	IDG 1110
	NGP(JNGP1) = NGP1	IDG 1120
160	CONTINUE	IDG 1130
	C DETERMINES GRID POINTS OUTSIDE THE DATA AREA.	IDG 1140
	C - IN SEMI-INFINITE RECTANGULAR AREA.	IDG 1150
	DO 450 IL0=1,NL0	IDG 1160
	NGP0 = 0	IDG 1170
	NGP1 = 0	IDG 1180
	IL0T3 = IL0*3	IDG 1190
	IP1 = IPL(IL0T3-2)	IDG 1200
	IP2 = IPL(IL0T3-1)	IDG 1210
	X1 = XD(IP1)	IDG 1220
	Y1 = YD(IP1)	IDG 1230
	X2 = XD(IP2)	IDG 1240
	Y2 = YD(IP2)	IDG 1250
	XMN = XIMN	IDG 1260
	XXM = XIMX	IDG 1270
	YMN = YIMN	IDG 1280
	YMX = YIMX	IDG 1290
	IF (Y2.GE.Y1) XMN = AMIN1(X1,X2)	IDG 1300
	IF (Y2.LE.Y1) XXM = AMAX1(X1,X2)	IDG 1310
	IF (X2.LE.X1) YMN = AMIN1(Y1,Y2)	IDG 1320
	IF (X2.GE.X1) YMX = AMAX1(Y1,Y2)	IDG 1330
	INSD = 0	IDG 1340
	DO 180 IXI=1,NXI0	IDG 1350
	IF (XI(IXI).GE.XMN .AND. XI(IXI).LE.XMX) GO TO 170	IDG 1360
	IF (INSD.EQ.0) GO TO 180	IDG 1370
	IXIMX = IXI - 1	IDG 1380
	GO TO 190	IDG 1390
170	IF (INSD.EQ.1) GO TO 180	IDG 1400
	INSD = 1	IDG 1410
	IXIMN = IXI	IDG 1420
180	CONTINUE	IDG 1430
	IF (INSD.EQ.0) GO TO 310	IDG 1440
	IXIMX = NXI0	IDG 1450
190	DO 300 IYI=1,NYI0	IDG 1460
	YII = YI(IYI)	IDG 1470
	IF (YII.LT.YMN .OR. YII.GT.YMX) GO TO 300	IDG 1480
	DO 290 IXI=IXIMN,IXIMX	IDG 1490
	XII = XI(IXI)	IDG 1500
	L = 0	IDG 1510
	IF (SIDE(X1,Y1,X2,Y2,XII,YII)) 210, 200, 290	IDG 1520
200	L = 1	IDG 1530
210	IF (SPDT(X2,Y2,X1,Y1,XII,YII)) 290, 220, 230	IDG 1540
220	L = 1	IDG 1550
230	IF (SPDT(X1,Y1,X2,Y2,XII,YII)) 290, 240, 250	IDG 1560
240	L = 1	IDG 1570
250	IZI = NXI0*(IYI-1) + IXI	IDG 1580
	IF (L.EQ.1) GO TO 260	IDG 1590
	NGP0 = NGP0 + 1	IDG 1600
	JIGP0 = JIGP0 + 1	IDG 1610
	IGP(JIGP0) = IZI	IDG 1620

	GO TO 290	IDG 1630
260	IF (JIGP1.GT.NXINYI) GO TO 280	IDG 1640
	DO 270 JIGP1I=JIGP1,NXINYI	IDG 1650
	IF (IZI.EQ.IGP(JIGP1I)) GO TO 290	IDG 1660
270	CONTINUE	IDG 1670
280	NGP1 = NGP1 + 1	IDG 1680
	JIGP1 = JIGP1 - 1	IDG 1690
	IGP(JIGP1) = IZI	IDG 1700
290	CONTINUE	IDG 1710
300	CONTINUE	IDG 1720
310	JNGP0 = JNGP0 + 1	IDG 1730
	NGP(JNGP0) = NGP0	IDG 1740
	JNGP1 = JNGP1 - 1	IDG 1750
	NGP(JNGP1) = NGP1	IDG 1760
	C - IN SEMI-INFINITE TRIANGULAR AREA.	IDG 1770
	NGP0 = 0	IDG 1780
	NGP1 = 0	IDG 1790
	ILP1 = MOD(IL0,NL0) + 1	IDG 1800
	ILP1T3 = ILP1*3	IDG 1810
	IP3 = IPL(ILP1T3-1)	IDG 1820
	X3 = XD(IP3)	IDG 1830
	Y3 = YD(IP3)	IDG 1840
	XMN = YIMN	IDG 1850
	XXM = XIMX	IDG 1860
	YMN = YIMN	IDG 1870
	YMX = YIMX	IDG 1880
	IF (Y3.GE.Y2 .AND. Y2.GE.Y1) XMN = X2	IDG 1890
	IF (Y3.LE.Y2 .AND. Y2.LE.Y1) XXM = X2	IDG 1900
	IF (X3.LE.X2 .AND. X2.LE.X1) YMN = Y2	IDG 1910
	IF (X3.GE.X2 .AND. X2.GE.X1) YMX = Y2	IDG 1920
	INSD = 0	IDG 1930
	DO 330 IXI=1,NXI0	IDG 1940
	IF (XI(IXI).GE.XMN .AND. XI(IXI).LE.XMX) GO TO 320	IDG 1950
	IF (INSD.EQ.0) GO TO 330	IDG 1960
	IXIMX = IXI - 1	IDG 1970
	GO TO 340	IDG 1980
320	IF (INSD.EQ.1) GO TO 330	IDG 1990
	INSD = 1	IDG 2000
	IXIMN = IXI	IDG 2010
330	CONTINUE	IDG 2020
	IF (INSD.EQ.0) GO TO 440	IDG 2030
	IXIMX = NXI0	IDG 2040
340	DO 430 IYI=1,NYI0	IDG 2050
	YII = YI(IYI)	IDG 2060
	IF (YII.LT.YMN .OR. YII.GT.YMX) GO TO 430	IDG 2070
	DO 420 IXI=IXIMN,IXIMX	IDG 2080
	XII = XI(IXI)	IDG 2090
	L = 0	IDG 2100
	IF (SPDT(X1,Y1,X2,Y2,XII,YII)) 360, 350, 420	IDG 2110
350	L = 1	IDG 2120
360	IF (SPDT(X3,Y3,X2,Y2,XII,YII)) 380, 370, 420	IDG 2130
370	L = 1	IDG 2140
380	IZI = NXI0*(IYI-1) + IXI	IDG 2150
	IF (L.EQ.1) GO TO 390	IDG 2160
	NGP0 = NGP0 + 1	IDG 2170
	JIGP0 = JIGP0 + 1	IDG 2180
	IGP(JIGP0) = IZI	IDG 2190
	GO TO 420	IDG 2200
390	IF (JIGP1.GT.NXINYI) GO TO 410	IDG 2210
	DO 400 JIGP1I=JIGP1,NXINYI	IDG 2220
	IF (IZI.EQ.IGP(JIGP1I)) GO TO 420	IDG 2230
400	CONTINUE	IDG 2240
410	NGP1 = NGP1 + 1	IDG 2250
	JIGP1 = JIGP1 - 1	IDG 2260
	IGP(JIGP1) = IZI	IDG 2270
420	CONTINUE	IDG 2280
430	CONTINUE	IDG 2290
440	JNGP0 = JNGP0 + 1	IDG 2300
	NGP(JNGP0) = NGP0	IDG 2310
	JNGP1 = JNGP1 - 1	IDG 2320
	NGP(JNGP1) = NGP1	IDG 2330
450	CONTINUE	IDG 2340
	RETURN	IDG 2350
	END	IDG 2360

```

SUBROUTINE IDLCTN(NDP, XD, YD, NT, IPT, NL, IPL, XII, YII, ITI, IDL 10
* IWK, WK) IDL 20
C THIS SUBROUTINE LOCATES A POINT, I.E., DETERMINES TO WHAT TRI- IDL 30
C ANGLE A GIVEN POINT (XII,YII) BELONGS. WHEN THE GIVEN POINT IDL 40
C DOES NOT LIE INSIDE THE DATA AREA, THIS SUBROUTINE DETERMINES IDL 50
C THE BORDER LINE SEGMENT WHEN THE POINT LIES IN AN OUTSIDE IDL 60
C RECTANGULAR AREA, AND TWO BORDER LINE SEGMENTS WHEN THE POINT IDL 70
C LIES IN AN OUTSIDE TRIANGULAR AREA. IDL 80
C THE INPUT PARAMETERS ARE IDL 90
C NDP = NUMBER OF DATA POINTS, IDL 100
C XD,YD = ARRAYS OF DIMENSION NDP CONTAINING THE X AND Y IDL 110
C COORDINATES OF THE DATA POINTS, IDL 120
C NT = NUMBER OF TRIANGLES, IDL 130
C IPT = INTEGER ARRAY OF DIMENSION 3*NT CONTAINING THE IDL 140
C POINT NUMBERS OF THE VERTEXES OF THE TRIANGLES, IDL 150
C NL = NUMBER OF BORDER LINE SEGMENTS, IDL 160
C IPL = INTEGER ARRAY OF DIMENSION 3*NL CONTAINING THE IDL 170
C POINT NUMBERS OF THE END POINTS OF THE BORDER IDL 180
C LINE SEGMENTS AND THEIR RESPECTIVE TRIANGLE IDL 190
C NUMBERS, IDL 200
C XII,YII = X AND Y COORDINATES OF THE POINT TO BE IDL 210
C LOCATED. IDL 220
C THE OUTPUT PARAMETER IS IDL 230
C ITI = TRIANGLE NUMBER, WHEN THE POINT IS INSIDE THE IDL 240
C DATA AREA, OR IDL 250
C TWO BORDER LINE SEGMENT NUMBERS, IL1 AND IL2, IDL 260
C CODED TO IL1*(NT+NL)+IL2, WHEN THE POINT IS IDL 270
C OUTSIDE THE DATA AREA. IDL 280
C THE OTHER PARAMETERS ARE IDL 290
C IWK = INTEGER ARRAY OF DIMENSION 18*NDP USED INTER- IDL 300
C NALLY AS A WORK AREA, IDL 310
C WK = ARRAY OF DIMENSION 8*NDP USED INTERNALLY AS A IDL 320
C WORK AREA. IDL 330
C DECLARATION STATEMENTS IDL 340
DIMENSION XD(100), YD(100), IPT(585), IPL(300), IWK(1800), IDL 350
* WK(800) IDL 360
DIMENSION NTSC(9), IDSC(9) IDL 370
COMMON /IDLC/ NIT IDL 380
C STATEMENT FUNCTIONS IDL 390
SIDE(U1,V1,U2,V2,U3,V3) = (U1-U3)*(V2-V3) - (V1-V3)*(U2-U3) IDL 400
SPDT(U1,V1,U2,V2,U3,V3) = (U1-U2)*(U3-U2) + (V1-V2)*(V3-V2) IDL 410
C PRELIMINARY PROCESSING IDL 420
NDP0 = NDP IDL 430
NT0 = NT IDL 440
NL0 = NL IDL 450
NTL = NT0 + NL0 IDL 460
X0 = XII IDL 470
Y0 = YII IDL 480
C PROCESSING FOR A NEW SET OF DATA POINTS IDL 490
IF (NIT.NE.0) GO TO 80 IDL 500
NIT = 1 IDL 510
C - DIVIDES THE X-Y PLANE INTO NINE RECTANGULAR SECTIONS. IDL 520
XMN = XD(1) IDL 530
XMX = XMN IDL 540
YMN = YD(1) IDL 550
YMX = YMN IDL 560
DO 10 IDP=2,NDP0 IDL 570
XI = XD(IDP) IDL 580
YI = YD(IDP) IDL 590
XMN = AMIN1(XI,XMN) IDL 600
XMX = AMAX1(XI,XMX) IDL 610
YMN = AMIN1(YI,YMN) IDL 620
YMX = AMAX1(YI,YMX) IDL 630
10 CONTINUE IDL 640
XS1 = (XMN+XMN+XMX)/3.0 IDL 650
XS2 = (XMN+XMX+XMX)/3.0 IDL 660
YS1 = (YMN+YMN+YMX)/3.0 IDL 670
YS2 = (YMN+YMX+YMX)/3.0 IDL 680
C - DETERMINES AND STORES IN THE IWK ARRAY TRIANGLE NUMBERS OF IDL 690
C - THE TRIANGLES ASSOCIATED WITH EACH OF THE NINE SECTIONS. IDL 700
DO 20 ISC=1,9 IDL 710
NTSC(ISC) = 0 IDL 720
IDSC(ISC) = 0 IDL 730
20 CONTINUE IDL 740
IT0T3 = 0 IDL 750
JWK = 0 IDL 760

```

```

DO 70 IT0=1,NT0
  IT0T3 = IT0T3 + 3
  I1 = IPT(IT0T3-2)
  I2 = IPT(IT0T3-1)
  I3 = IPT(IT0T3)
  XMN = AMIN1(XD(I1),XD(I2),XD(I3))
  XMX = AMAX1(XD(I1),XD(I2),XD(I3))
  YMN = AMIN1(YD(I1),YD(I2),YD(I3))
  YMX = AMAX1(YD(I1),YD(I2),YD(I3))
  IF (YMN.GT.YS1) GO TO 30
  IF (XMN.LE.XS1) IDSC(1) = 1
  IF (XMX.GE.XS1 .AND. XMN.LE.XS2) IDSC(2) = 1
  IF (XMX.GE.XS2) IDSC(3) = 1
30 IF (YMX.LT.YS1 .OR. YMN.GT.YS2) GO TO 40
  IF (XMN.LE.XS1) IDSC(4) = 1
  IF (XMX.GE.XS1 .AND. XMN.LE.XS2) IDSC(5) = 1
  IF (XMX.GE.XS2) IDSC(6) = 1
40 IF (YMX.LT.YS2) GO TO 50
  IF (XMN.LE.XS1) IDSC(7) = 1
  IF (XMX.GE.XS1 .AND. XMN.LE.XS2) IDSC(8) = 1
  IF (XMX.GE.XS2) IDSC(9) = 1
50 DO 60 ISC=1,9
  IF (IDSC(ISC).EQ.0) GO TO 60
  JIWK = 9*NTSC(ISC) + ISC
  IWK(JIWK) = IT0
  NTSC(ISC) = NTSC(ISC) + 1
  IDSC(ISC) = 0
60 CONTINUE
C - STORES IN THE WK ARRAY THE MINIMUM AND MAXIMUM OF THE X AND
C - Y COORDINATE VALUES FOR EACH OF THE TRIANGLE.
  JWK = JWK + 4
  WK(JWK-3) = XMN
  WK(JWK-2) = XMX
  WK(JWK-1) = YMN
  WK(JWK) = YMX
70 CONTINUE
GO TO 110
C CHECKS IF IN THE SAME TRIANGLE AS PREVIOUS.
80 IT0 = ITIPV
  IF (IT0.GT.NT0) GO TO 90
  IT0T3 = IT0*3
  IP1 = IPT(IT0T3-2)
  X1 = XD(IP1)
  Y1 = YD(IP1)
  IP2 = IPT(IT0T3-1)
  X2 = XD(IP2)
  Y2 = YD(IP2)
  IF (SIDE(X1,Y1,X2,Y2,X0,Y0).LT.0.0) GO TO 110
  IP3 = IPT(IT0T3)
  X3 = XD(IP3)
  Y3 = YD(IP3)
  IF (SIDE(X2,Y2,X3,Y3,X0,Y0).LT.0.0) GO TO 110
  IF (SIDE(X3,Y3,X1,Y1,X0,Y0).LT.0.0) GO TO 110
  GO TO 170
C CHECKS IF ON THE SAME BORDER LINE SEGMENT.
90 IL1 = IT0/NTL
  IL2 = IT0 - IL1*NTL
  IL1T3 = IL1*3
  IP1 = IPL(IL1T3-2)
  X1 = XD(IP1)
  Y1 = YD(IP1)
  IP2 = IPL(IL1T3-1)
  X2 = XD(IP2)
  Y2 = YD(IP2)
  IF (IL2.NE.IL1) GO TO 100
  IF (SPDT(X1,Y1,X2,Y2,X0,Y0).LT.0.0) GO TO 110
  IF (SPDT(X2,Y2,X1,Y1,X0,Y0).LT.0.0) GO TO 110
  IF (SIDE(X1,Y1,X2,Y2,X0,Y0).GT.0.0) GO TO 110
  GO TO 170
C CHECKS IF BETWEEN THE SAME TWO BORDER LINE SEGMENTS.
100 IF (SPDT(X1,Y1,X2,Y2,X0,Y0).GT.0.0) GO TO 110
  IP3 = IPL(3*IL2-1)
  X3 = XD(IP3)
  Y3 = YD(IP3)
  IF (SPDT(X3,Y3,X2,Y2,X0,Y0).LE.0.0) GO TO 170
C LOCATES INSIDE THE DATA AREA.

```

```

IDL 770
IDL 780
IDL 790
IDL 800
IDL 810
IDL 820
IDL 830
IDL 840
IDL 850
IDL 860
IDL 870
IDL 880
IDL 890
IDL 900
IDL 910
IDL 920
IDL 930
IDL 940
IDL 950
IDL 960
IDL 970
IDL 980
IDL 990
IDL 1000
IDL 1010
IDL 1020
IDL 1030
IDL 1040
IDL 1050
IDL 1060
IDL 1070
IDL 1080
IDL 1090
IDL 1100
IDL 1110
IDL 1120
IDL 1130
IDL 1140
IDL 1150
IDL 1160
IDL 1170
IDL 1180
IDL 1190
IDL 1200
IDL 1210
IDL 1220
IDL 1230
IDL 1240
IDL 1250
IDL 1260
IDL 1270
IDL 1280
IDL 1290
IDL 1300
IDL 1310
IDL 1320
IDL 1330
IDL 1340
IDL 1350
IDL 1360
IDL 1370
IDL 1380
IDL 1390
IDL 1400
IDL 1410
IDL 1420
IDL 1430
IDL 1440
IDL 1450
IDL 1460
IDL 1470
IDL 1480
IDL 1490
IDL 1500
IDL 1510
IDL 1520

```

```

C - DETERMINES THE SECTION IN WHICH THE POINT IN QUESTION LIES.
110 ISC = 1
    IF (X0.GE.XS1) ISC = ISC + 1
    IF (X0.GE.XS2) ISC = ISC + 1
    IF (Y0.GE.YS1) ISC = ISC + 3
    IF (Y0.GE.YS2) ISC = ISC + 3
C - SEARCHES THROUGH THE TRIANGLES ASSOCIATED WITH THE SECTION.
NTSCI = NTSC(ISC)
IF (NTSCI.LE.0) GO TO 130
JIWK = -9 + ISC
DO 120 ITSC=1,NTSCI
    JIWK = JIWK + 9
    IT0 = IWK(JIWK)
    JWK = IT0*4
    IF (X0.LT.WK(JWK-3)) GO TO 120
    IF (X0.GT.WK(JWK-2)) GO TO 120
    IF (Y0.LT.WK(JWK-1)) GO TO 120
    IF (Y0.GT.WK(JWK)) GO TO 120
    IT0T3 = IT0*3
    IP1 = IPT(IT0T3-2)
    X1 = XD(IP1)
    Y1 = YD(IP1)
    IP2 = IPT(IT0T3-1)
    X2 = XD(IP2)
    Y2 = YD(IP2)
    IF (SIDE(X1,Y1,X2,Y2,X0,Y0).LT.0.0) GO TO 120
    IP3 = IPT(IT0T3)
    X3 = XD(IP3)
    Y3 = YD(IP3)
    IF (SIDE(X2,Y2,X3,Y3,X0,Y0).LT.0.0) GO TO 120
    IF (SIDE(X3,Y3,X1,Y1,X0,Y0).LT.0.0) GO TO 120
    GO TO 170
120 CONTINUE
C LOCATES OUTSIDE THE DATA AREA.
130 DO 150 IL1=1,NL0
    IL1T3 = IL1*3
    IP1 = IPL(IL1T3-2)
    X1 = XD(IP1)
    Y1 = YD(IP1)
    IP2 = IPL(IL1T3-1)
    X2 = XD(IP2)
    Y2 = YD(IP2)
    IF (SPDT(X2,Y2,X1,Y1,X0,Y0).LT.0.0) GO TO 150
    IF (SPDT(X1,Y1,X2,Y2,X0,Y0).LT.0.0) GO TO 140
    IF (SIDE(X1,Y1,X2,Y2,X0,Y0).GT.0.0) GO TO 150
    IL2 = IL1
    GO TO 160
140 IL2 = MOD(IL1,NL0) + 1
    IP3 = IPL(3*IL2-1)
    X3 = XD(IP3)
    Y3 = YD(IP3)
    IF (SPDT(X3,Y3,X2,Y2,X0,Y0).LE.0.0) GO TO 160
150 CONTINUE
    IT0 = 1
    GO TO 170
160 IT0 = IL1*NTL + IL2
C NORMAL EXIT
170 ITI = IT0
    ITIPV = IT0
    RETURN
    END

```

SUBROUTINE IDPDRV(NDP,XD,YD,ZD,NCP,IPC,PD)

```

C THIS SUBROUTINE ESTIMATES PARTIAL DERIVATIVES OF THE FIRST AND
C SECOND ORDER AT THE DATA POINTS.
C THE INPUT PARAMETERS ARE
C   NDI = NUMBER OF DATA POINTS,
C   XD,YD,ZD = ARRAYS OF DIMENSION NDP CONTAINING THE X,
C             Y, AND Z COORDINATES OF THE DATA POINTS,
C   NCP = NUMBER OF ADDITIONAL DATA POINTS USED FOR ESTI-
C         MATING PARTIAL DERIVATIVES AT EACH DATA POINT,
C   IPC = INTEGER ARRAY OF DIMENSION NCP*NDP CONTAINING
C         THE POINT NUMBERS OF NCP DATA POINTS CLOSEST TO
C         EACH OF THE NDP DATA POINTS.
C THE OUTPUT PARAMETER IS

```

IDL 1530
IDL 1540
IDL 1550
IDL 1560
IDL 1570
IDL 1580
IDL 1590
IDL 1600
IDL 1610
IDL 1620
IDL 1630
IDL 1640
IDL 1650
IDL 1660
IDL 1670
IDL 1680
IDL 1690
IDL 1700
IDL 1710
IDL 1720
IDL 1730
IDL 1740
IDL 1750
IDL 1760
IDL 1770
IDL 1780
IDL 1790
IDL 1800
IDL 1810
IDL 1820
IDL 1830
IDL 1840
IDL 1850
IDL 1860
IDL 1870
IDL 1880
IDL 1890
IDL 1900
IDL 1910
IDL 1920
IDL 1930
IDL 1940
IDL 1950
IDL 1960
IDL 1970
IDL 1980
IDL 1990
IDL 2000
IDL 2010
IDL 2020
IDL 2030
IDL 2040
IDL 2050
IDL 2060
IDL 2070
IDL 2080
IDL 2090
IDL 2100
IDL 2110
IDL 2120
IDL 2130

ID008940
ID008950
ID008960
ID008970
ID008980
ID008990
ID009000
ID009010
ID009020
ID009030
ID009040
ID009050
ID009060

```

C   PD = ARRAY OF DIMENSION 5*NDP, WHERE THE ESTIMATED
C   ZX, ZY, ZXX, ZXY, AND ZYY VALUES AT THE DATA
C   POINTS ARE TO BE STORED.
C DECLARATION STATEMENTS
      DIMENSION  XD(100), YD(100), ZD(100), IPC(400), PD(500)
      REAL        NMX, NMY, NMZ, NMXX, NMXY, NMYX, NMYY
C PRELIMINARY PROCESSING
  10 NDP0=NDP
      NCP0=NCP
      NCPM1=NCP0-1
C ESTIMATION OF ZX AND ZY
  20 DO 24 IP0=1, NDP0
      X0=XD(IP0)
      Y0=YD(IP0)
      Z0=ZD(IP0)
      NMX=0.0
      NMY=0.0
      NMZ=0.0
      JIPC0=NCP0*(IP0-1)
      DO 23 IC1=1, NCPM1
          JIPC=JIPC0+IC1
          IPI=IPC(JIPC)
          DX1=XD(IPI)-X0
          DY1=YD(IPI)-Y0
          DZ1=ZD(IPI)-Z0
          IC2MN=IC1+1
          DO 22 IC2=IC2MN, NCP0
              JIPC=JIPC0+IC2
              IPI=IPC(JIPC)
              DX2=XD(IPI)-X0
              DY2=YD(IPI)-Y0
              DNMZ=DX1*DY2-DY1*DX2
              IF (DNMZ.EQ.0.0) GO TO 22
              DZ2=ZD(IPI)-Z0
              DNMX=DY1*DZ2-DZ1*DY2
              DNMY=DZ1*DX2-DX1*DZ2
              IF (DNMZ.GE.0.0) GO TO 21
              DNMX=-DNMX
              DNMY=-DNMY
              DNMZ=-DNMZ
  21      NMX=NMX+DNMX
          NMY=NMY+DNMY
          NMZ=NMZ+DNMZ
  22      CONTINUE
  23      CONTINUE
          JPD0=5*IP0
          PD(JPD0-4)=-NMX/NMZ
          PD(JPD0-3)=-NMY/NMZ
  24      CONTINUE
C ESTIMATION OF ZXX, ZXY, AND ZYY
  30 DO 34 IP0=1, NDP0
      JPD0=JPD0+5
      X0=XD(IP0)
      JPD0=5*IP0
      Y0=YD(IP0)
      ZX0=PD(JPD0-4)
      ZY0=PD(JPD0-3)
      NMXX=0.0
      NMXY=0.0
      NMYX=0.0
      NMYY=0.0
      NMZ =0.0
      JIPC0=NCP0*(IP0-1)
      DO 33 IC1=1, NCPM1
          JIPC=JIPC0+IC1
          IPI=IPC(JIPC)
          DX1=XD(IPI)-X0
          DY1=YD(IPI)-Y0
          JPD=5*IPI
          DZX1=PD(JPD-4)-ZX0
          DZY1=PD(JPD-3)-ZY0
          IC2MN=IC1+1
          DO 32 IC2=IC2MN, NCP0
              JIPC=JIPC0+IC2
              IPI=IPC(JIPC)
              DX2=XD(IPI)-X0

```

```

ID009070
ID009080
ID009090
ID009100
ID009110
ID009120
ID009130
ID009140
ID009150
ID009160
ID009170
ID009180
ID009190
ID009200
ID009210
ID009220
ID009230
ID009240
ID009250
ID009260
ID009270
ID009280
ID009290
ID009300
ID009310
ID009320
ID009330
ID009340
ID009350
ID009360
ID009370
ID009380
ID009390
ID009400
ID009410
ID009420
ID009430
ID009440
ID009450
ID009460
ID009470
ID009480
ID009490
ID009500
ID009510
ID009520
ID009530
ID009540
ID009550
ID009560
ID009570
ID009580
ID009590
ID009600
ID009610
ID009620
ID009630
ID009640
ID009650
ID009660
ID009670
ID009680
ID009690
ID009700
ID009710
ID009720
ID009730
ID009740
ID009750
ID009760
ID009770
ID009780
ID009790
ID009800
ID009810
ID009820

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          DY2=YD(IPI)-Y0
          DNMZ =DX1*DY2 -DY1*DX2
          IF(DNMZ.EQ.0.0) GO TO 32
          JPD=5*IPI
          DZX2=PD(JPD-4)-ZX0
          DZY2=PD(JPD-3)-ZY0
          DNMXX=DY1*DZX2-DZX1*DY2
          DNMXY=DZX1*DX2-DX1*DZX2
          DNMYY=DY1*DZY2-DZY1*DY2
          DNMZY=DZY1*DX2-DX1*DZY2
          IF(DNMZY.GE.0.0) GO TO 31
          DNMXX=-DNMXX
          DNMXY=-DNMXY
          DNMYY=-DNMYY
          DNMZ =-DNMZ
31      NMXX=NMXX+DNMXX
          NMXY=NMXY+DNMXY
          NMYY=NMYY+DNMYY
          NMZ =NMZ +DNMZ
32      CONTINUE
33      CONTINUE
          PD(JPD-2)=-NMXX/NMZ
          PD(JPD-1)=-NMXY+NMYY/(2.0*NMZ)
          PD(JPD) =-NMYY/NMZ
34      CONTINUE
          RETURN
          END
          ID009830
          ID009840
          ID009850
          ID009860
          ID009870
          ID009880
          ID009890
          ID009900
          ID009910
          ID009920
          ID009930
          ID009940
          ID009950
          ID009960
          ID009970
          ID009980
          ID009990
          ID010000
          ID010010
          ID010020
          ID010030
          ID010040
          ID010050
          ID010060
          ID010070
          ID010080
          ID010090
          ID010100
          ID010110

          SUBROUTINE IDPTIP(XD,YD,ZD,NT,IPT,NL,IPL,PDD,ITI,XII,YII,
1              ZII)
          ID010190
          ID010200
C THIS SUBROUTINE PERFORMS PUNCTUAL INTERPOLATION OR EXTRAPOLA-
C TION, I.E., DETERMINES THE Z VALUE AT A POINT.
          ID010210
          ID010220
C THE INPUT PARAMETERS ARE
          ID010230
C XD,YD,ZD = ARRAYS OF DIMENSION NDP CONTAINING THE X,
C Y, AND Z COORDINATES OF THE DATA POINTS, WHERE
          ID010240
          ID010250
C NDP IS THE NUMBER OF THE DATA POINTS,
          ID010260
C NT = NUMBER OF TRIANGLES,
          ID010270
C IPT = INTEGER ARRAY OF DIMENSION 3*NT CONTAINING THE
C POINT NUMBERS OF THE VERTEXES OF THE TRIANGLES,
          ID010280
          ID010290
C NL = NUMBER OF BORDER LINE SEGMENTS,
          ID010300
C IPL = INTEGER ARRAY OF DIMENSION 3*NL CONTAINING THE
C POINT NUMBERS OF THE END POINTS OF THE BORDER
          ID010310
          ID010320
C LINE SEGMENTS AND THEIR RESPECTIVE TRIANGLE
          ID010330
C NUMBERS,
          ID010340
C PDD = APPAY OF DIMENSION 5*NDP CONTAINING THE PARTIAL
C DERIVATIVES AT THE DATA POINTS,
          ID010350
          ID010360
C ITI = TRIANGLE NUMBER OF THE TRIANGLE IN WHICH LIES
          ID010370
C THE POINT FOR WHICH INTERPOLATION IS TO BE
          ID010380
C PERFORMED,
          ID010390
C XII,Y*II = X AND Y COORDINATES OF THE POINT FOR WHICH
          ID010400
C INTERPOLATION IS TO BE PERFORMED.
          ID010410
C THE OUTPUT PARAMETER IS
          ID010420
C ZII = INTERPOLATED Z VALUE.
          ID010430
C DECLARATION STATEMENTS
          ID010440
          DIMENSION XD(100),YD(100),ZD(100),IPT(585),IPL(300),
          ID010450
          1 PDD(500)
          ID010460
          COMMON/IDPI/ITPV
          ID010470
          DIMENSION X(3),Y(3),Z(3),PD(15),
          ID010480
          1 ZU(3),ZV(3),ZUU(3),ZUV(3),ZVV(3)
          ID010490
          REAL LU,LV
          ID010500
          EQUIVALENCE (P5,P50)
          ID010510
C PRELIMINARY PROCESSING
          ID010520
          10 IT0=ITI
          ID010530
          NTL=NT+NL
          ID010540
          IF(IT0.LE.NTL) GO TO 20
          ID010550
          IL1=IT0/NTL
          ID010560
          IL2=IT0-IL1*NTL
          ID010570
          IF(17.1.EQ.IL2) GO TO 40
          ID010580
          GO TO 60
          ID010590
C CALCULATION OF ZII BY INTERPOLATION.
          ID010600
C CHECKS IF THE NECESSARY COEFFICIENTS HAVE BEEN CALCULATED.
          ID010610
          20 IF(IT0.EQ.ITPV) GO TO 30
          ID010620
C LOADS COORDINATE AND PARTIAL DERIVATIVE VALUES AT THE
          ID010630

```

```

C VERTEXES.
21 JIPT=3*(IT0-1)
   JPD=0
   DO 23 I=1,3
     JIPT=JIPT+1
     IDP=IPT(JIPT)
     X(I)=XD(IDP)
     Y(I)=YD(IDP)
     Z(I)=ZD(IDP)
     JPDD=5*(IDP-1)
     DO 22 KPD=1,5
       JPD=JPD+1
       JPDD=JPDD+1
       PD(JPD)=PDD(JPDD)
     22 CONTINUE
   23 CONTINUE
C DETERMINES THE COEFFICIENTS FOR THE COORDINATE SYSTEM
C TRANSFORMATION FROM THE X-Y SYSTEM TO THE U-V SYSTEM
C AND VICE VERSA.
24 X0=X(1)
   Y0=Y(1)
   A=X(2)-X0
   B=X(3)-X0
   C=Y(2)-Y0
   D=Y(3)-Y0
   AD=A*D
   BC=B*C
   DLT=AD-BC
   AP= D/DLT
   BP=-B/DLT
   CP=-C/DLT
   DP= A/DLT
C CONVERTS THE PARTIAL DERIVATIVES AT THE VERTEXES OF THE
C TRIANGLE FOR THE U-V COORDINATE SYSTEM.
25 AA=A*A
   ACT2=2.0*A*C
   CC=C*C
   AB=A*B
   ADBC=AD+BC
   CD=C*D
   BB=B*B
   BDT2=2.0*B*D
   DD=D*D
DO 26 I=1,3
   IPD=5*I
   ZU(I)=A*PD(JPD-4)+C*PD(JPD-3)
   ZV(I)=B*PD(JPD-4)+D*PD(JPD-3)
   ZUU(I)=AA*PD(JPD-2)+ACT2*PD(JPD-1)+CC*PD(JPD)
   ZUV(I)=AB*PD(JPD-2)+ADBC*PD(JPD-1)+CD*PD(JPD)
   ZVV(I)=BB*PD(JPD-2)+BDT2*PD(JPD-1)+DD*PD(JPD)
26 CONTINUE
C CALCULATES THE COEFFICIENTS OF THE POLYNOMIAL.
27 P00=Z(1)
   P10=ZU(1)
   P01=ZV(1)
   P20=0.5*ZUU(1)
   P11=ZUV(1)
   P02=0.5*ZVV(1)
   H1=Z(2)-P00-P10-P20
   H2=ZU(2)-P10-ZUU(1)
   H3=ZUU(2)-ZUU(1)
   P30= 10.0*H1-4.0*H2+0.5*H3
   P40=-15.0*H1+7.0*H2 -H3
   P50= 6.0*H1-3.0*H2+0.5*H3
   H1=Z(3)-P00-P01-P02
   H2=ZV(3)-P01-ZVV(1)
   H3=ZVV(3)-ZVV(1)
   P03= 10.0*H1-4.0*H2+0.5*H3
   P04=-15.0*H1+7.0*H2 -H3
   P05= 6.0*H1-3.0*H2+0.5*H3
   LU=SQRT(AA+CC)
   LV=SQRT(BB+DD)
   THXU=ATAN2(C,A)
   THUV=ATAN2(D,B)-THXU
   CSUV=COS(THUV)
   P41=5.0*LV*CSUV/LU*P50
ID010640
ID010650
ID010660
ID010670
ID010680
ID010690
ID010700
ID010710
ID010720
ID010730
ID010740
ID010750
ID010760
ID010770
ID010780
ID010790
ID010800
ID010810
ID010820
ID010830
ID010840
ID010850
ID010860
ID010870
ID010880
ID010890
ID010900
ID010910
ID010920
ID010930
ID010940
ID010950
ID010960
ID010970
ID010980
ID010990
ID011000
ID011010
ID011020
ID011030
ID011040
ID011050
ID011060
ID011070
ID011080
ID011090
ID011100
ID011110
ID011120
ID011130
ID011140
ID011150
ID011160
ID011170
ID011180
ID011190
ID011200
ID011210
ID011220
ID011230
ID011240
ID011250
ID011260
ID011270
ID011280
ID011290
ID011300
ID011310
ID011320
ID011330
ID011340
ID011350
ID011360
ID011370
ID011380
ID011390

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P14=5.0*LU*CSUV/LV*P05
H1=ZV(2)-P01-P11-P41
H2=ZUV(2)-P11-4.0*P41
P21= 3.0*H1-H2
P31=-2.0*H1+H2
H1=ZU(3)-P10-P11-P14
H2=ZUV(3)-P11-4.0*P14
P12= 3.0*H1-H2
P13=-2.0*H1+H2
THUS=ATAN2(D-C,B-A)-THXU
THSV=THUV-THUS
AA= SIN(THSV)/LU
BB=-COS(THSV)/LU
CC= SIN(THUS)/LV
DD= COS(THUS)/LV
AC=AA*CC
AD=AA*DD
BC=BB*CC
G1=AA*AC*(3.0*BC+2.0*AD)
G2=CC*AC*(3.0*AD+2.0*BC)
H1=-AA*AA*AA*(5.0*AA*BB*P50+(4.0*BC+AD)*P41)
1 -CC*CC*CC*(5.0*CC*DD*P05+(4.0*AD+BC)*P14)
H2=0.5*ZVV(2)-P02-P12
H3=0.5*ZUU(3)-P20-P21
P22=(G1*H2+G2*H3-H1)/(G1+G2)
P32=H2-P22
P23=H3-P22
ITPV=IT0
C CONVERTS XII AND YII TO U-V SYSTEM.
30 DX=XII-X0
DY=YII-Y0
U=AP*DX+BP*DY
V=CP*DX+DP*DY
C EVALUATES THE POLYNOMIAL.
31 P0=P00+V*(P01+V*(P02+V*(P03+V*(P04+V*P05))))
P1=P10+V*(P11+V*(P12+V*(P13+V*P14)))
P2=P20+V*(P21+V*(P22+V*P23))
P3=P30+V*(P31+V*P32)
P4=P40+V*P41
ZII=P0+U*(P1+U*(P2+U*(P3+U*(P4+U*P5)))
RETURN
C CALCULATION OF ZII BY EXTRAPOLATION IN THE RECTANGLE.
C CHECKS IF THE NECESSARY COEFFICIENTS HAVE BEEN CALCULATED.
40 IF(IT0,EQ,ITPV) GO TO 50
C LOADS COORDINATE AND PARTIAL DERIVATIVE VALUES AT THE END
C POINTS OF THE BORDER LINE SEGMENT.
41 JIPL=3*(IL1-1)
JPD=0
DO 43 I=1,2
JIPL=JIPL+1
IDP=IPL(JIPL)
X(I)=XD(IDP)
Y(I)=YD(IDP)
Z(I)=ZD(IDP)
JPDD=5*(IDP-1)
DO 42 KPD=1,5
JPD=JPD+1
JPDD=JPDD+1
PD(JPD)=PDD(JPDD)
42 CONTINUE
43 CONTINUE
C DETERMINES THE COEFFICIENTS FOR THE COORDINATE SYSTEM
C TRANSFORMATION FROM THE X-Y SYSTEM TO THE U-V SYSTEM
C AND VICE VERSA.
44 X0=X(1)
Y0=Y(1)
A=Y(2)-Y(1)
B=X(2)-X(1)
C=-B
D=A
AD=A*D
BC=B*C
DLT=AD-BC
AP= D/DLT
BP=-B/DLT
CP=-BP
ID011400
ID011410
ID011420
ID011430
ID011440
ID011450
ID011460
ID011470
ID011480
ID011490
ID011500
ID011510
ID011520
ID011530
ID011540
ID011550
ID011560
ID011570
ID011580
ID011590
ID011600
ID011610
ID011620
ID011630
ID011640
ID011650
ID011660
ID011670
ID011680
ID011690
ID011700
ID011710
ID011720
ID011730
ID011740
ID011750
ID011760
ID011770
ID011780
ID011790
ID011800
ID011810
ID011820
ID011830
ID011840
ID011850
ID011860
ID011870
ID011880
ID011890
ID011900
ID011910
ID011920
ID011930
ID011940
ID011950
ID011960
ID011970
ID011980
ID011990
ID012000
ID012010
ID012020
ID012030
ID012040
ID012050
ID012060
ID012070
ID012080
ID012090
ID012100
ID012110
ID012120
ID012130
ID012140
ID012150

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DP= AP
C CONVERTS THE PARTIAL DERIVATIVES AT THE END POINTS OF THE
C BORDER LINE SEGMENT FOR THE U-V COORDINATE SYSTEM.
45 AA=A*A
ACT2=2.0*A*C
CC=C*C
AB=A*B
ADBC=AD+BC
CD=C*D
BB=B*B
BDT2=2.0*B*D
DD=D*D
DO 46 I=1,2
JPD=5*I
ZU(I)=A*PD(JPD-4)+C*PD(JPD-3)
ZV(I)=B*PD(JPD-4)+D*PD(JPD-3)
ZUU(I)=AA*PD(JPD-2)+ACT2*PD(JPD-1)+CC*PD(JPD)
ZUV(I)=AB*PD(JPD-2)+ADBC*PD(JPD-1)+CD*PD(JPD)
ZVV(I)=BB*PD(JPD-2)+BDT2*PD(JPD-1)+DD*PD(JPD)
46 CONTINUE
C CALCULATES THE COEFFICIENTS OF THE POLYNOMIAL.
47 P00=Z(1)
P10=ZU(1)
P01=ZV(1)
P20=0.5*ZUU(1)
P11=ZUV(1)
P02=0.5*ZVV(1)
H1=Z(2)-P00-P01-P02
H2=ZV(2)-P01-ZVV(1)
H3=ZVV(2)-ZVV(1)
P03= 10.0*H1-4.0*H2+0.5*H3
P04=-15.0*H1+7.0*H2 -H3
P05= 6.0*H1-3.0*H2+0.5*H3
H1=ZU(2)-P10-P11
H2=ZUV(2)-P11
P12= 3.0*H1-H2
P13=-2.0*H1+H2
P21=0.0
P23=-ZUU(2)+ZUU(1)
P22=-1.5*P23
ITPV=IT0
C CONVERTS XII AND YII TO U-V SYSTEM.
50 DX=XII-X0
DY=YII-Y0
U=AP*DX+BP*DY
V=CP*DX+DP*DY
C EVALUATES THE POLYNOMIAL.
51 P0=P00+V*(P01+V*(P02+V*(P03+V*(P04+V*P05))))
P1=P10+V*(P11+V*(P12+V*P13))
P2=P20+V*(P21+V*(P22+V*P23))
ZII=P0+U*(P1+U*P2)
RETURN
C CALCULATION OF ZII BY EXTRAPOLATION IN THE TRIANGLE.
C CHECKS IF THE NECESSARY COEFFICIENTS HAVE BEEN CALCULATED.
60 IF(IT0.EQ.ITPV) GO TO 70
C LOADS COORDINATE AND PARTIAL DERIVATIVE VALUES AT THE VERTEX
C OF THE TRIANGLE.
61 JIPL=3*IL2-2
IDP=IPL(JIPL)
X(1)=XD(IDP)
Y(1)=YD(IDP)
Z(1)=ZD(IDP)
JPDD=5*(IDP-1)
DO 62 KPD=1,5
JPDD=JPDD+1
PD(KPD)=PDD(JPDD)
62 CONTINUE
C CALCULATES THE COEFFICIENTS OF THE POLYNOMIAL.
67 P00=Z(1)
P10=PD(1)
P01=PD(2)
P20=0.5*PD(3)
P11=PD(4)
P02=0.5*PD(5)
ITPV=IT0
C CONVERTS XII AND YII TO U-V SYSTEM.

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ID012160
ID012170
ID012180
ID012190
ID012200
ID012210
ID012220
ID012230
ID012240
ID012250
ID012260
ID012270
ID012280
ID012290
ID012300
ID012310
ID012320
ID012330
ID012340
ID012350
ID012360
ID012370
ID012380
ID012390
ID012400
ID012410
ID012420
ID012430
ID012440
ID012450
ID012460
ID012470
ID012480
ID012490
ID012500
ID012510
ID012520
ID012530
ID012540
ID012550
ID012560
ID012570
ID012580
ID012590
ID012600
ID012610
ID012620
ID012630
ID012640
ID012650
ID012660
ID012670
ID012680
ID012690
ID012700
ID012710
ID012720
ID012730
ID012740
ID012750
ID012760
ID012770
ID012780
ID012790
ID012800
ID012810
ID012820
ID012830
ID012840
ID012850
ID012860
ID012870
ID012880
ID012890
ID012900
ID012910

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```

70 U=XII-X(1)
V=YII-Y(1)
C EVALUATES THE POLYNOMIAL.
71 P0=P00+V*(P01+V*P02)
P1=P10+V*P11
ZII=P0+U*(P1+U*P20)
RETURN
END
ID012920
ID012930
ID012940
ID012950
ID012960
ID012970
ID012980
ID012990

SUBROUTINE IDSFFT(MD,NCP,NDP,XD,YD,ZD,NXI,NYI,XI,YI,ZI,
1 IWK,WK)
ID013070
ID013080
C THIS SUBROUTINE PERFORMS SMOOTH SURFACE FITTING WHEN THE PRO-
C JECTIONS OF THE DATA POINTS IN THE X-Y PLANE ARE IRREGULARLY
C DISTRIBUTED IN THE PLANE.
ID013090
ID013100
C THE INPUT PARAMETERS ARE
ID013110
ID013120
C MD = MODE OF COMPUTATION (MUST BE 1, 2, OR 3),
ID013130
C = 1 FOR NEW NCP AND/OR NEW XD-YD,
ID013140
C = 2 FOR OLD NCP, OLD XD-YD, NEW XI-YI,
ID013150
C = 3 FOR OLD NCP, OLD XD-YD, OLD XI-YI,
ID013160
C NCP = NUMBER OF ADDITIONAL DATA POINTS USED FOR ESTI-
ID013170
C MATING PARTIAL DERIVATIVES AT EACH DATA POINT
ID013180
C (MUST BE 2 OR GREATER, BUT SMALLER THAN NDP),
ID013190
C NDP = NUMBER OF DATA POINTS (MUST BE 4 OR GREATER),
ID013200
C XD = ARRAY OF DIMENSION NDP CONTAINING THE X
ID013210
C COORDINATES OF THE DATA POINTS,
ID013220
C YD = ARRAY OF DIMENSION NDP CONTAINING THE Y
ID013230
C COORDINATES OF THE DATA POINTS,
ID013240
C ZD = ARRAY OF DIMENSION NDP CONTAINING THE Z
ID013250
C COORDINATES OF THE DATA POINTS,
ID013260
C NXI = NUMBER OF OUTPUT GRID POINTS IN THE X COORDINATE
ID013270
C (MUST BE 1 OR GREATER),
ID013280
C NYI = NUMBER OF OUTPUT GRID POINTS IN THE Y COORDINATE
ID013290
C (MUST BE 1 OR GREATER),
ID013300
C XI = ARRAY OF DIMENSION NXI CONTAINING THE X
ID013310
C COORDINATES OF THE OUTPUT GRID POINTS,
ID013320
C YI = ARRAY OF DIMENSION NYI CONTAINING THE Y
ID013330
C COORDINATES OF THE OUTPUT GRID POINTS.
ID013340
C THE OUTPUT PARAMETER IS
ID013350
C ZI = DOUBLY-DIMENSIONED ARRAY OF DIMENSION (NXI,NYI),
ID013360
C WHERE THE INTERPOLATED Z VALUES AT THE OUTPUT
ID013370
C GRID POINTS ARE TO BE STORED.
ID013380
C THE OTHER PARAMETERS ARE
ID013390
C IWK = INTEGER ARRAY OF DIMENSION
ID013400
C MAX0(31,27+NCP)*NDP+NXI*NYI
ID013410
C USED INTERNALLY AS A WORK AREA,
ID013420
C WK = ARRAY OF DIMENSION 5*NDP USED INTERNALLY AS A
ID013430
C WORK AREA.
ID013440
C THE VERY FIRST CALL TO THIS SUBROUTINE AND THE CALL WITH A NEW
ID013450
C NCP VALUE, A NEW NDP VALUE, AND/OR NEW CONTENTS OF THE XD AND
ID013460
C YD ARRAYS MUST BE MADE WITH MD=1. THE CALL WITH MD=2 MUST BE
ID013470
C PRECEDED BY ANOTHER CALL WITH THE SAME NCP AND NDP VALUES AND
ID013480
C WITH THE SAME CONTENTS OF THE XD AND YD ARRAYS. THE CALL WITH
ID013490
C MD=3 MUST BE PRECEDED BY ANOTHER CALL WITH THE SAME NCP, NDP,
ID013500
C NXI, AND NYI VALUES AND WITH THE SAME CONTENTS OF THE XD, YD,
ID013510
C XI, AND YI ARRAYS. BETWEEN THE CALL WITH MD=2 OR MD=3 AND ITS
ID013520
C PRECEDING CALL, THE IWK AND WK ARRAYS MUST NOT BE DISTURBED.
ID013530
C USE OF A VALUE BETWEEN 3 AND 5 (INCLUSIVE) FOR NCP IS RECOM-
ID013540
C MENDED UNLESS THERE ARE EVIDENCES THAT DICTATE OTHERWISE.
ID013550
C THE LUN CONSTANT IN THE DATA INITIALIZATION STATEMENT IS THE
ID013560
C LOGICAL UNIT NUMBER OF THE STANDARD OUTPUT UNIT AND IS,
ID013570
C THEREFORE, SYSTEM DEPENDENT.
ID013580
C THIS SUBROUTINE CALLS THE IDCLDP, IDGRID, IDPDRV, IDPTIP, AND
ID013590
C IDTANG SUBROUTINES.
ID013600
C DECLARATION STATEMENTS
ID013610
C DIMENSION XD(100),YD(100),ZD(100),XI(101),YI(101),
ID013620
C 1 ZI(10201),IWK(13301),WK(500)
ID013630
C COMMON/IDPI/ITPV
ID013640
C DATA LUN/6/
ID013650
C SETTING OF SOME INPUT PARAMETERS TO LOCAL VARIABLES.
ID013660
C (FOR MD=1,2,3)
ID013670
10 MD0=MD
ID013680
NCP0=NCP
ID013690
NDP0=NDP
ID013700
NXI0=NXI
ID013710
NYI0=NYI
ID013720

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C ERROR CHECK. (FOR MD=1,2,3)
20 IF (MD0.LT.1.OR.MD0.GT.3) GO TO 90
   IF (NCP0.LT.2.OR.NCP0.GE.NDP0) GO TO 90
   IF (NDP0.LT.4) GO TO 90
   IF (NXI0.LT.1.OR.NYI0.LT.1) GO TO 90
   IF (MD0.GE.2) GO TO 21
   IWK(1)=NCP0
   IWK(2)=NDP0
   GO TO 22
21 NCPPV=IWK(1)
   NDPPV=IWK(2)
   IF (NCP0.NE.NCPPV) GO TO 90
   IF (NDP0.NE.NDPPV) GO TO 90
22 IF (MD0.GE.3) GO TO 23
   IWK(3)=NXI0
   IWK(4)=NYI0
   GO TO 30
23 NXIPV=IWK(3)
   NYIPV=IWK(4)
   IF (NXI0.NE.NXIPV) GO TO 90
   IF (NYI0.NE.NYIPV) GO TO 90
C ALLOCATION OF STORAGE AREAS IN THE IWK ARRAY. (FOR MD=1,2,3)
30 JWIPT=16
   JWIWL=6*NDP0+1
   JWNGP0=JWIWL-1
   JWIPL=24*NDP0+1
   JWIWP=30*NDP0+1
   JWIPC=27*NDP0+1
   JWIGP0=MAX0(31,27+NCP0)*NDP0
C TRIANGULATES THE X-Y PLANE. (FOR MD=1)
40 IF (MD0.GT.1) GO TO 50
   CALL IDTANG (NDP0,XD,YD,NT,IWK(JWIPT),NL,IWK(JWIPL),
1     IWK(JWIWL),IWK(JWIWP),WK)
   IWK(5)=NT
   IWK(6)=NL
   IF (NT.EQ.0) RETURN
C DETERMINES NCP POINTS CLOSEST TO EACH DATA POINT. (FOR MD=1)
50 IF (MD0.GT.1) GO TO 60
   CALL IDCLDP (NDP0,XD,YD,NCP0,IWK(JWIPC))
   IF (IWK(JWIPC).EQ.0) RETURN
C SORTS OUTPUT GRID POINTS IN ASCENDING ORDER OF THE TRIANGLE
C NUMBER AND THE BORDER LINE SEGMENT NUMBER. (FOR MD=1,2)
60 IF (MD0.EQ.3) GO TO 70
   CALL IDGRID (XD,YD,NT,IWK(JWIPT),NL,IWK(JWIPL),NXI0,NYI0,
1     XI,YI,IWK(JWNGP0+1),IWK(JWIGP0+1))
C ESTIMATES PARTIAL DERIVATIVES AT ALL DATA POINTS.
C (FOR MD=1,2,3)
70 CALL IDPRV (NDP0,XD,YD,ZD,NCP0,IWK(JWIPC),WK)
C INTERPOLATES THE ZI VALUES. (FOR MD=1,2,3)
80 ITPV=0
   JIG0MX=0
   JIG0MN=NXI0*NYI0+1
   NNGP=NT+2*NL
   DO 89 JNGP=1,NNGP
     ITI=JNGP
     IF (JNGP.LE.NT) GO TO 81
     IL1=(JNGP-NT+1)/2
     IL2=(JNGP-NT+2)/2
     IF (IL2.GT.NL) IL2=1
     ITI=IL1*(NT+NL)+IL2
81 JWNGP=JWNGP0+JNGP
   NGP0=IWK(JWNGP)
   IF (NGP0.EQ.0) GO TO 86
   JIG0MN=JIG0MX+1
   JIG0MX=JIG0MX+NGP0
   DO 82 JIGP=JIG0MN,JIG0MX
     JWIGP=JWIGP0+JIGP
     IZI=IWK(JWIGP)
     IYI=(IZI-1)/NXI0+1
     IXI=IZI-NXI0*(IYI-1)
     CALL IDPTIP (XD,YD,ZD,NT,IWK(JWIPT),NL,IWK(JWIPL),WK,
1     ITI,XI(IXI),YI(IYI),ZI(IZI))
82 CONTINUE
86 JWNGP=JWNGP0+2*NNGP+1-JNGP
   NGP1=IWK(JWNGP)
   IF (NGP1.EQ.0) GO TO 89

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ID013730
ID013740
ID013750
ID013760
ID013770
ID013780
ID013790
ID013800
ID013810
ID013820
ID013830
ID013840
ID013850
ID013860
ID013870
ID013880
ID013890
ID013900
ID013910
ID013920
ID013930
ID013940
ID013950
ID013960
ID013970
ID013980
ID013990
ID014000
ID014010
ID014020
ID014030
ID014040
ID014050
ID014060
ID014070
ID014080
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ID014100
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ID014470
ID014480

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JIGIMX=JIGIMN-1
JIGIMN=JIGIMN-NGP1
DO 87 JIGP=JIGIMN,JIGIMX
  JWIGP=JWIGP+JIGP
  IZI=IWK(JWIGP)
  IYI=(IZI-1)/NXI+1
  IXI=IZI-NXI*(IYI-1)
  CALL IDPTIP(XD,YD,ZD,NT,IWK(JWIPT),NL,IWK(JWIPL),WK,
1      ITI,XI(IXI),YI(IYI),ZI(IZI))
87 CONTINUE
89 CONTINUE
  RETURN
C ERROR EXIT
  90 WRITE (LUN,2090) MD0,NCP0,NDP0,NXI0,NYI0
  RETURN
C FORMAT STATEMENT FOR ERROR MESSAGE
2090 FORMAT(1X/41H *** IMPROPER INPUT PARAMETER VALUE(S) ./
1 7H MD =,I4,10X,5HNCP =,I6,10X,5HNDP =,I6,
2 10X,5HNXI =,I6,10X,5HNYI =,I6/
3 35H ERROR DETECTED IN ROUTINE IDSFFT/)
END

SUBROUTINE IDTANG(NDP,XD,YD,NT,IPT,NL,IPL,IWL,IWP,WK)
C THIS SUBROUTINE PERFORMS TRIANGULATION. IT DIVIDES THE X-Y
C PLANE INTO A NUMBER OF TRIANGLES ACCORDING TO GIVEN DATA
C POINTS IN THE PLANE, DETERMINES LINE SEGMENTS THAT FORM THE
C BORDER OF DATA AREA, AND DETERMINES THE TRIANGLE NUMBERS
C CORRESPONDING TO THE BORDER LINE SEGMENTS.
C AT COMPLETION, POINT NUMBERS OF THE VERTEXES OF EACH TRIANGLE
C ARE LISTED COUNTER-CLOCKWISE. POINT NUMBERS OF THE END POINTS
C OF EACH BORDER LINE SEGMENT ARE LISTED COUNTER-CLOCKWISE,
C LISTING ORDER OF THE LINE SEGMENTS BEING COUNTER-CLOCKWISE.
C THE LUN CONSTANT IN THE DATA INITIALIZATION STATEMENT IS THE
C LOGICAL UNIT NUMBER OF THE STANDARD OUTPUT UNIT AND IS,
C THEREFORE, SYSTEM DEPENDENT.
C THIS SUBROUTINE CALLS THE IDXCHG FUNCTION.
C THE INPUT PARAMETERS ARE
C NDP = NUMBER OF DATA POINTS,
C XD = ARRAY OF DIMENSION NDP CONTAINING THE
C X COORDINATES OF THE DATA POINTS,
C YD = ARRAY OF DIMENSION NDP CONTAINING THE
C Y COORDINATES OF THE DATA POINTS.
C THE OUTPUT PARAMETERS ARE
C NT = NUMBER OF TRIANGLES,
C IPT = INTEGER ARRAY OF DIMENSION 6*NDP-15, WHERE THE
C POINT NUMBERS OF THE VERTEXES OF THE (IT)TH
C TRIANGLE ARE TO BE STORED AS THE (3*IT-2)ND,
C (3*IT-1)ST, AND (3*IT)TH ELEMENTS,
C IT=1,2,...,NT,
C NL = NUMBER OF BORDER LINE SEGMENTS,
C IPL = INTEGER ARRAY OF DIMENSION 6*NDP, WHERE THE
C POINT NUMBERS OF THE END POINTS OF THE (IL)TH
C BORDER LINE SEGMENT AND ITS RESPECTIVE TRIANGLE
C NUMBER ARE TO BE STORED AS THE (3*IL-2)ND,
C (3*IL-1)ST, AND (3*IL)TH ELEMENTS,
C IL=1,2,...,NL.
C THE OTHER PARAMETERS ARE
C IWL = INTEGER ARRAY OF DIMENSION 18*NDP USED
C INTERNALLY AS A WORK AREA,
C IWP = INTEGER ARRAY OF DIMENSION NDP USED
C INTERNALLY AS A WORK AREA,
C WK = ARRAY OF DIMENSION NDP USED INTERNALLY AS A
C WORK AREA.
C DECLARATION STATEMENTS
  DIMENSION XD(100),YD(100),IPT(585),IPL(600),
1 IWL(1800),IWP(100),WK(100)
  DIMENSION ITF(2)
  DATA RATIO/1.0E-6/,NREF/100/,LUN/6/
C STATEMENT FUNCTIONS
  DSQF(U1,V1,U2,V2)=(U2-U1)**2+(V2-V1)**2
  SIDE(U1,V1,U2,V2,U3,V3)=(V3-V1)*(U2-U1)-(U3-U1)*(V2-V1)
C PRELIMINARY PROCESSING
10 NDP=NDP
  NDPM1=NDP-1
  IF(NDP.LT.4) GO TO 90
C DETERMINES THE CLOSEST PAIR OF DATA POINTS AND THEIR MIDPOINT.
20 DSQMN=DSQF(XD(1),YD(1),XD(2),YD(2))

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ID015310

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IPMN1=1
IPMN2=2
DO 22 IP1=1,NDPM1
  X1=XD(IP1)
  Y1=YD(IP1)
  IP1P1=IP1+1
  DO 21 IP2=IP1P1,NDP0
    DSQI=DSQF(X1,Y1,XD(IP2),YD(IP2))
    IF(DSQI.EQ.0.0) GO TO 91
    IF(DSQI.GE.DSQMN) GO TO 21
    DSQMN=DSQI
    IPMN1=IP1
    IPMN2=IP2
21 CONTINUE
22 CONTINUE
DSQ12=DSQMN
XDMP=(XD(IPMN1)+XD(IPMN2))/2.0
YDMP=(YD(IPMN1)+YD(IPMN2))/2.0
C SORTS THE OTHER (NDP-2) DATA POINTS IN ASCENDING ORDER OF
C DISTANCE FROM THE MIDPOINT AND STORES THE SORTED DATA POINT
C NUMBERS IN THE IWP ARRAY.
30 JP1=2
DO 31 IP1=1,NDP0
  IF(IP1.EQ.IPMN1.OR.IP1.EQ.IPMN2) GO TO 31
  JP1=JP1+1
  IWP(JP1)=IP1
  WK(JP1)=DSQF(XDMP,YDMP,XD(IP1),YD(IP1))
31 CONTINUE
DO 33 JP1=3,NDPM1
  DSQMN=WK(JP1)
  JPMN=JP1
  DO 32 JP2=JP1,NDP0
    IF(WK(JP2).GE.DSQMN) GO TO 32
    DSQMN=WK(JP2)
    JPMN=JP2
32 CONTINUE
ITS=IWP(JP1)
IWP(JP1)=IWP(JPMN)
IWP(JPMN)=ITS
WK(JPMN)=WK(JP1)
33 CONTINUE
C IF NECESSARY, MODIFIES THE ORDERING IN SUCH A WAY THAT THE
C FIRST THREE DATA POINTS ARE NOT COLLINEAR.
35 AR=DSQ12*RATIO
X1=XD(IPMN1)
Y1=YD(IPMN1)
DX21=XD(IPMN2)-X1
DY21=YD(IPMN2)-Y1
DO 36 JP=3,NDP0
  IP=IWP(JP)
  IF(ABS((YD(IP)-Y1)*DX21-(XD(IP)-X1)*DY21).GT.AR)
1 GO TO 37
36 CONTINUE
GO TO 92
37 IF(JP.EQ.3) GO TO 40
JPMX=JP
JP=JPMX+1
DO 38 JPC=4,JPMX
  JP=JP-1
  IWP(JP)=IWP(JP-1)
38 CONTINUE
IWP(3)=IP
C FORMS THE FIRST TRIANGLE. STORES POINT NUMBERS OF THE VER-
C TEXES OF THE TRIANGLE IN THE IPT ARRAY, AND STORES POINT NUM-
C BERS OF THE BORDER LINE SEGMENTS AND THE TRIANGLE NUMBER IN
C THE IPL ARRAY.
40 IP1=IPMN1
IP2=IPMN2
IP3=IWP(3)
IF(SIDE(XD(IP1),YD(IP1),XD(IP2),YD(IP2),XD(IP3),YD(IP3))
1 .GE.0.0) GO TO 41
IP1=IPMN2
IP2=IPMN1
41 NT0=1
NTT3=3
IPT(1)=IP1

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IPT(2)=IP2
IPT(3)=IP3
NLØ=3
NLT3=9
IPL(1)=IP1
IPL(2)=IP2
IPL(3)=1
IPL(4)=IP2
IPL(5)=IP3
IPL(6)=1
IPL(7)=IP3
IPL(8)=IP1
IPL(9)=1
C ADDS THE REMAINING (NDP-3) DATA POINTS, ONE BY ONE.
5Ø DO 79 JP1=4,NDPØ
    IP1=IWP(JP1)
    X1=XD(IP1)
    Y1=YD(IP1)
C - DETERMINES THE VISIBLE BORDER LINE SEGMENTS.
    IP2=IPL(1)
    JPMN=1
    DXMN=XD(IP2)-X1
    DYMN=YD(IP2)-Y1
    DSQMN=DXMN**2+DYMN**2
    ARMN=DSQMN*RATIO
    JPMX=1
    DXMX=DXMN
    DYMx=DYMN
    DSQMX=DSQMN
    ARMX=ARMN
DO 52 JP2=2,NLØ
    IP2=IPL(3*JP2-2)
    DX=XD(IP2)-X1
    DY=YD(IP2)-Y1
    AR=DY*DXMN-DX*DYMN
    IF(AR.GT.ARMN) GO TO 51
    DSQI=DX**2+DY**2
    IF(AR.GE.(-ARMN).AND.DSQI.GE.DSQMN) GO TO 51
    JPMN=JP2
    DXMN=DX
    DYMN=DY
    DSQMN=DSQI
    ARMN=DSQMN*RATIO
51 AR=DY*DXMX-DX*DYMx
    IF(AR.LT.(-ARMx)) GO TO 52
    DSQI=DX**2+DY**2
    IF(AR.LE.ARMx.AND.DSQI.GE.DSQMX) GO TO 52
    JP2=JP2
    DXMX=DX
    DYMx=DY
    DSQMX=DSQI
    ARMX=DSQMX*RATIO
52 CONTINUE
    IF(JPMx.LT.JPMN) JPMx=JPMx+NLØ
    NSH=JPMN-1
    IF(NSH.LE.Ø) GO TO 6Ø
C - SHIFTS (ROTATES) THE IPL ARRAY TO HAVE THE INVISIBLE BORDER
C - LINE SEGMENTS CONTAINED IN THE FIRST PART OF THE IPL ARRAY.
    NSHT3=NSH*3
DO 53 JP2T3=3,NSHT3,3
    JP3T3=JP2T3+NLT3
    IPL(JP3T3-2)=IPL(JP2T3-2)
    IPL(JP3T3-1)=IPL(JP2T3-1)
    IPL(JP3T3) =IPL(JP2T3)
53 CONTINUE
DO 54 JP2T3=3,NLT3,3
    JP3T3=JP2T3+NSHT3
    IPL(JP2T3-2)=IPL(JP3T3-2)
    IPL(JP2T3-1)=IPL(JP3T3-1)
    IPL(JP2T3) =IPL(JP3T3)
54 CONTINUE
    JPMx=JPMx-NSH
C - ADDS TRIANGLES TO THE IPT ARRAY, UPDATES BORDER LINE
C - SEGMENTS IN THE IPL ARRAY, AND SETS FLAGS FOR THE BORDER
C - LINE SEGMENTS TO BE REEXAMINED IN THE IWL ARRAY.
6Ø JWL=Ø

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DO 64 JP2=JPMX,NLØ
  JP2T3=JP2*3
  IPL1=IPL(JP2T3-2)
  IPL2=IPL(JP2T3-1)
  IT =IPL(JP2T3)
C - - ADDS A TRIANGLE TO THE IPT ARRAY.
  NTØ=NTØ+1
  NTT3=NTT3+3
  IPT(NTT3-2)=IPL2
  IPT(NTT3-1)=IPL1
  IPT(NTT3) =IP1
C - - UPDATES BORDER LINE SEGMENTS IN THE IPL ARRAY.
  IF(JP2.NE.JPMX) GO TO 61
  IPL(JP2T3-1)=IP1
  IPL(JP2T3) =NTØ
61 IF(JP2.NE.NLØ) GO TO 62
  NLN=JPMX+1
  NLNT3=NLN*3
  IPL(NLNT3-2)=IP1
  IPL(NLNT3-1)=IPL(1)
  IPL(NLNT3) =NTØ
C - - DETERMINES THE VERTEX THAT DOES NOT LIE ON THE BORDER
C - - LINE SEGMENTS.
62 ITT3=IT*3
  IPTI=IPT(ITT3-2)
  IF(IPTI.NE.IPL1.AND.IPTI.NE.IPL2) GO TO 63
  IPTI=IPT(ITT3-1)
  IF(IPTI.NE.IPL1.AND.IPTI.NE.IPL2) GO TO 63
  IPTI=IPT(ITT3)
C - - CHECKS IF THE EXCHANGE IS NECESSARY.
63 IF(IXCHG(XD,YD,IP1,IPTI,IPL1,IPL2).EQ.Ø) GO TO 64
C - - MODIFIES THE IPT ARRAY WHEN NECESSARY.
  IPT(ITT3-2)=IPTI
  IPT(ITT3-1)=IPL1
  IPT(ITT3) =IP1
  IPT(NTT3-1)=IPTI
  IF(JP2.EQ.JPMX) IPL(JP2T3)=IT
  IF(JP2.EQ.NLØ.AND.IPL(3).EQ.IT) IPL(3)=NTØ
C - - SETS FLAGS IN THE IWL ARRAY.
  JWL=JWL+4
  IWL(JWL-3)=IPL1
  IWL(JWL-2)=IPTI
  IWL(JWL-1)=IPTI
  IWL(JWL) =IPL2
64 CONTINUE
  NLØ=NLN
  NLT3=NLNT3
  NLF=JWL/2
  IF(NLF.EQ.Ø) GO TO 79
C - IMPROVES TRIANGULATION.
7Ø NTT3P3=NTT3+3
  DO 78 IREP=1,NREP
    DO 76 ILF=1,NLF
      ILFT2=ILF*2
      IPL1=IWL(ILFT2-1)
      IPL2=IWL(ILFT2)
C - - LOCATES IN THE IPT ARRAY TWO TRIANGLES ON BOTH SIDES OF
C - - THE FLAGGED LINE SEGMENT.
      NTF=Ø
      DO 71 ITT3R=3,NTT3,3
        ITT3=NTT3P3-ITT3R
        IPT1=IPT(ITT3-2)
        IPT2=IPT(ITT3-1)
        IPT3=IPT(ITT3)
        IF(IPL1.NE.IPT1.AND.IPL1.NE.IPT2.AND.
1          IPL1.NE.IPT3) GO TO 71
        IF(IPL2.NE.IPT1.AND.IPL2.NE.IPT2.AND.
1          IPL2.NE.IPT3) GO TO 71
        NTF=NTF+1
        ITF(NTF)=ITT3/3
        IF(NTF.EQ.2) GO TO 72
71 CONTINUE
        IF(NTF.LT.2) GO TO 76
C - - DETERMINES THE VERTEXES OF THE TRIANGLES THAT DO NOT LIE
C - - ON THE LINE SEGMENT.
72 ITT3=ITF(1)*3

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IPTI1=IPT(IT1T3-2) ID017600
IF (IPTI1.NE.IPL1.AND.IPTI1.NE.IPL2) GO TO 73 ID017610
IPTI1=IPT(IT1T3-1) ID017620
IF (IPTI1.NE.IPL1.AND.IPTI1.NE.IPL2) GO TO 73 ID017630
IPTI1=IPT(IT1T3) ID017640
73 IT2T3=ITF(2)*3 ID017650
IPTI2=IPT(IT2T3-2) ID017660
IF (IPTI2.NE.IPL1.AND.IPTI2.NE.IPL2) GO TO 74 ID017670
IPTI2=IPT(IT2T3-1) ID017680
IF (IPTI2.NE.IPL1.AND.IPTI2.NE.IPL2) GO TO 74 ID017690
IPTI2=IPT(IT2T3) ID017700
C -- CHECKS IF THE EXCHANGE IS NECESSARY. ID017710
74 IF (IDXCHG(XD,YD,IPTI1,IPTI2,IPL1,IPL2).EQ.0) ID017720
1 GO TO 76 ID017730
C -- MODIFIES THE IPT ARRAY WHEN NECESSARY. ID017740
IPT(IT1T3-2)=IPTI1 ID017750
IPT(IT1T3-1)=IPTI2 ID017760
IPT(IT1T3) =IPL1 ID017770
IPT(IT2T3-2)=IPTI2 ID017780
IPT(IT2T3-1)=IPTI1 ID017790
IPT(IT2T3) =IPL2 ID017800
C -- SETS NEW FLAGS. ID017810
JWL=JWL+8 ID017820
IWL(JWL-7)=IPL1 ID017830
IWL(JWL-6)=IPTI1 ID017840
IWL(JWL-5)=IPTI1 ID017850
IWL(JWL-4)=IPL2 ID017860
IWL(JWL-3)=IPL2 ID017870
IWL(JWL-2)=IPTI2 ID017880
IWL(JWL-1)=IPTI2 ID017890
IWL(JWL) =IPL1 ID017900
DO 75 JLT3=3,NLT3,3 ID017910
IPLJ1=IPL(JLT3-2) ID017920
IPLJ2=IPL(JLT3-1) ID017930
IF((IPLJ1.EQ.IPL1.AND.IPLJ2.EQ.IPTI2).OR. ID017940
1 (IPLJ2.EQ.IPL1.AND.IPLJ1.EQ.IPTI2)) ID017950
2 IPL(JLT3)=ITF(1) ID017960
IF((IPLJ1.EQ.IPL2.AND.IPLJ2.EQ.IPTI1).OR. ID017970
1 (IPLJ2.EQ.IPL2.AND.IPLJ1.EQ.IPTI1)) ID017980
2 IPL(JLT3)=ITF(2) ID017990
75 CONTINUE ID018000
76 CONTINUE ID018010
NLFC=NLF ID018020
NLF=JWL/2 ID018030
IF(NLF.EQ.NLFC) GO TO 79 ID018040
C -- RESETS THE IWL ARRAY FOR THE NEXT ROUND. ID018050
JWL=0 ID018060
JWL1MN=(NLFC+1)*2 ID018070
NLFT2=NLF*2 ID018080
DO 77 JWL1=JWL1MN,NLFT2,2 ID018090
JWL=JWL+2 ID018100
IWL(JWL-1)=IWL(JWL1-1) ID018110
IWL(JWL) =IWL(JWL1) ID018120
77 CONTINUE ID018130
NLF=JWL/2 ID018140
78 CONTINUE ID018150
79 CONTINUE ID018160
C REARRANGES THE IPT ARRAY SO THAT THE VERTEXES OF EACH TRIANGLE ID018170
C ARE LISTED COUNTER-CLOCKWISE. ID018180
80 DO 81 ITT3=3,NTT3,3 ID018190
IP1=IPT(ITT3-2) ID018200
IP2=IPT(ITT3-1) ID018210
IP3=IPT(ITT3) ID018220
IF(SIDE(XD(IP1),YD(IP1),XD(IP2),YD(IP2),XD(IP3),YD(IP3)) ID018230
1 .GE.0.0) GO TO 81 ID018240
IPT(ITT3-2)=IP2 ID018250
IPT(ITT3-1)=IP1 ID018260
81 CONTINUE ID018270
NT=NT0 ID018280
NL=NL0 ID018290
RETURN ID018300
C ERROR EXIT ID018310
90 WRITE (LUN,2090) NDP0 ID018320
GO TO 93 ID018330
91 WRITE (LUN,2091) NDP0,IP1,IP2,X1,Y1 ID018340
GO TO 93 ID018350

```

```

92 WRITE (LUN,2092) NDP0
93 WRITE (LUN,2093)
  NT=0
  RETURN
C FORMAT STATEMENTS
2090 FORMAT(1X/23H *** NDP LESS THAN 4./8H NDP =,I5)
2091 FORMAT(1X/29H *** IDENTICAL DATA POINTS./
  1 8H NDP =,I5,5X,5HIP1 =,I5,5X,5HIP2 =,I5,
  2 5X,4HXD =,E12.4,5X,4HYD =,E12.4)
2092 FORMAT(1X/33H *** ALL COLLINEAR DATA POINTS./
  1 8H NDP =,I5)
2093 FORMAT(35H ERROR DETECTED IN ROUTINE IDTANG/)
  END
ID018360
ID018370
ID018380
ID018390
ID018400
ID018410
ID018420
ID018430
ID018440
ID018450
ID018460
ID018470
ID018480

FUNCTION IDXCHG(X,Y,I1,I2,I3,I4)
C THIS FUNCTION DETERMINES WHETHER OR NOT THE EXCHANGE OF TWO
C TRIANGLES IS NECESSARY ON THE BASIS OF MAX-MIN-ANGLE CRITERION
C BY C. L. LAWSON.
C THE INPUT PARAMETERS ARE
C X,Y = ARRAYS CONTAINING THE COORDINATES OF THE DATA
C POINTS,
C I1,I2,I3,I4 = POINT NUMBERS OF FOUR POINTS P1, P2,
C P3, AND P4 THAT FORM A QUADRILATERAL WITH P3
C AND P4 CONNECTED DIAGONALLY.
C THIS FUNCTION RETURNS AN INTEGER VALUE 1 (ONE) WHEN AN EX-
C CHANGE IS NECESSARY, AND 0 (ZERO) OTHERWISE.
C DECLARATION STATEMENTS
  DIMENSION X(100),Y(100)
  EQUIVALENCE (C2SQ,C1SQ),(A3SQ,B2SQ),(B3SQ,A1SQ),
  1 (A4SQ,B1SQ),(B4SQ,A2SQ),(C4SQ,C3SQ)
C PRELIMINARY PROCESSING
  10 X1=X(I1)
  Y1=Y(I1)
  X2=X(I2)
  Y2=Y(I2)
  X3=X(I3)
  Y3=Y(I3)
  X4=X(I4)
  Y4=Y(I4)
C CALCULATION
  20 IDX=0
  U3=(Y2-Y3)*(X1-X3)-(X2-X3)*(Y1-Y3)
  U4=(Y1-Y4)*(X2-X4)-(X1-X4)*(Y2-Y4)
  IF(U3*U4.LE.0.0) GO TO 30
  U1=(Y3-Y1)*(X4-X1)-(X3-X1)*(Y4-Y1)
  U2=(Y4-Y2)*(X3-X2)-(X4-X2)*(Y3-Y2)
  A1SQ=(X1-X3)**2+(Y1-Y3)**2
  B1SQ=(X4-X1)**2+(Y4-Y1)**2
  C1SQ=(X3-X4)**2+(Y3-Y4)**2
  A2SQ=(X2-X4)**2+(Y2-Y4)**2
  B2SQ=(X3-X2)**2+(Y3-Y2)**2
  C3SQ=(X2-X1)**2+(Y2-Y1)**2
  S1SQ=U1*U1/(C1SQ*AMAX1(A1SQ,B1SQ))
  S2SQ=U2*U2/(C2SQ*AMAX1(A2SQ,B2SQ))
  S3SQ=U3*U3/(C3SQ*AMAX1(A3SQ,B3SQ))
  S4SQ=U4*U4/(C4SQ*AMAX1(A4SQ,B4SQ))
  IF(AMIN1(S1SQ,S2SQ).LT.AMIN1(S3SQ,S4SQ)) IDX=1
  30 IDXCHG=IDX
  RETURN
  END
ID018560
ID018570
ID018580
ID018590
ID018600
ID018610
ID018620
ID018630
ID018640
ID018650
ID018660
ID018670
ID018680
ID018690
ID018700
ID018710
ID018720
ID018730
ID018740
ID018750
ID018760
ID018770
ID018780
ID018790
ID018800
ID018810
ID018820
ID018830
ID018840
ID018850
ID018860
ID018870
ID018880
ID018890
ID018900
ID018910
ID018920
ID018930
ID018940
ID018950
ID018960
ID018970
ID018980
ID018990
ID019000
ID019010

```


*** IDSFFT ***

```
COMMON/IDPI/ITPV,DMMY(27)          ID013640
40 IF(MD0.GT.1) GO TO 41            ID014030
GO TO 50                             ID014081
41 NT=IWK(5)                         ID014082
NL=IWK(6)                             ID014083
```

After these changes are made, the algorithm works properly with the "reset own" option on the B6700 computer.

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ALGORITHM 527

A Fortran Implementation of the Generalized Marching Algorithm [D3]

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Key Words and Phrases: marching algorithms, block tridiagonal, elliptic partial differential equations
CR Categories: 5.14, 5.17
Language: Fortran

1. DESCRIPTION

Subroutines GMA and GMAS are implementations of the generalized marching algorithm [1, 2], which may be used to solve linear systems arising from 5-point discretizations of separable or constant coefficient elliptic boundary-value problems on rectangular domains. A Dirichlet, Neumann, or mixed boundary condition may be independently specified on each side of the rectangle; periodic boundary conditions may be specified on opposing sides. An appropriate linear system is block tridiagonal, or nearly block tridiagonal, real, symmetric, positive or negative definite (for GMA) or semidefinite (for GMAS), and of the form $Ax = b$:

$$\begin{bmatrix} T + \alpha_1 I & -\beta_2 I & & -\beta_1 I \\ -\beta_2 I & T + \alpha_2 I & -\beta_3 I & \\ & \ddots & \ddots & \\ & -\beta_{N-1} I & T + \alpha_{N-1} I & -\beta_N I \\ -\beta_1 I & & -\beta_N I & T + \alpha_N I \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_{N-1} \\ x_N \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_{N-1} \\ b_N \end{bmatrix} \quad (1)$$

Here $T = TM$ is the (nearly) tridiagonal $M \times M$ matrix

$$TM = \begin{bmatrix} \gamma_1 & -\sigma_2 & & & -\sigma_r \\ -\sigma_2 & \gamma_2 & -\sigma_3 & & \\ & \ddots & \ddots & \ddots & \\ & & -\sigma_{M-1} & \gamma_{M-1} & -\sigma_M \\ -\sigma_1 & & & -\sigma_M & \gamma_M \end{bmatrix} \quad (2)$$

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This program was written while the author was a summer visitor at Argonne National Laboratory, Argonne, IL.

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The α_i , β_i , γ_i , and σ_i are scalars, and x_i and b_i are M-vectors. Many common discretizations which do not automatically yield symmetric matrices A (for example, in the case of Neumann or mixed boundary conditions) can be transformed into the form of eq. (1) using diagonal similarity transformations [4, 11]. No restrictions on N and M, other than $N, M \geq 3$ are imposed.

The eigenvalues of certain principal submatrices of the $N \times N$ matrix TN, given by

$$\text{TN} = \begin{bmatrix} \alpha_1 & -\beta & & & & & & & -\beta_1 \\ -\beta_2 & \alpha_{22} & & & & & & & \\ & & \cdot & & & & & & \\ & & & \cdot & & & & & \\ & & & & \cdot & & & & \\ & & & & & \cdot & & & \\ & & & & & & \cdot & & \\ & & & & & & & \cdot & \\ & & & & & & & & \cdot \\ & & & & & & & & & -\beta_N \\ -\beta_1 & & & & & & & & & \alpha_N \end{bmatrix} \quad (3)$$

are required by the generalized marching algorithm and must be computed in a separate preprocessing phase. To solve a problem, the user must call GMA or GMAS twice, first to carry out the preprocessing, and then to solve the linear system. (If the same linear system is to be solved with many right-hand sides, the preprocessing need be done only once if the contents of the array R are saved.) Nineteen other subroutines are called internally as required.

PARTN, ROOTSC, ROOTSG, SORT, QL, and BANDR are entered during the preprocessing phase. In PARTN, the matrix A is partitioned into $ND + 1$ blocks by $ND = [N/K]$ separating lines (the M unknowns in the vector x_j comprise the J-th line). Each block has K or fewer lines, where K is the marching parameter specified by the user or supplied by default. The integer NE, which satisfies $2^{NE-1} \leq ND \leq 2^{NE} - 1$ if $ND > 0$, $NE = 0$ if $ND = 0$, is also computed.

ROOTSC and ROOTSG perform eigenvalue calculations for constant coefficient and general separable problems, respectively. SORT orders the eigenvalues according to size, both to insure numerical stability and to allow certain reductions in computation to occur in the constant coefficient case. QL and BANDR are called by ROOTSC and ROOTSG in connection with solving the eigenvalue problems; they are adaptations of TQLRAT and BANDR from the EISPACK subroutine library [7, 12].

In the backsolution phase, subroutines STEP1, STEP2, STEP3, STEP4, TRI1, TRI2, MARCH1, MARCH2, and MARCH3 are entered. In STEP1 the $ND + 1$ problems associated with the partitioned matrix A are partially solved using the marching algorithm. Next, in STEP2, $NE - 1$ "2-reductions" are applied to a reduced set of equations involving the separating lines. In STEP3, the solution vectors x_j corresponding to the separating lines are calculated in NE recursive substeps, effectively decomposing the N line problem into $ND + 1$ smaller problems. In STEP4, these smaller problems are solved via the marching algorithm. If periodic boundary conditions are specified for the matrix TN, certain adjustments are made in the backsolution process.

In the backsolution phase, $O(N \cdot NE)$ linear systems of the form $(T - rI)x = y$ are solved; this is done in TRI1 and TRI2. Additionally, the backsolution requires $O(N)$ matrix multiplications of the form $x = Ty$, which are carried out in MARCH1, MARCH2, and MARCH3.

Subroutine GMA is entered using

CALL GMA(NPC, N, NTYPE, TN1, TN2, M, MTYPE, TM1, TM2, IBDIM, B, K, R, A)

The parameter list is:

NPC is an integer. If $NPC = 0$, GMA does preprocessing calculations; if $NPC = 1$

GMA solves the linear system.

N is the dimension of the matrix TN; $N > 2$.

NTYPE is an integer describing the matrix TN.

NTYPE = 1: General separable; Dirichlet, Neumann, or mixed boundary conditions.

$$\beta_1 = 0; \beta_i \neq 0, 2 \leq i \leq N;$$

$$\alpha_i \text{ arbitrary, } 1 \leq i \leq N.$$

NTYPE = 2: Constant coefficient; Dirichlet, Neumann, or mixed boundary conditions.

$$\beta_1 = 0; \beta_i = \beta \neq 0, 3 \leq i \leq N - 1;$$

$$\alpha_i = \alpha, \alpha \text{ arbitrary, } 2 \leq i \leq N - 1.$$

Top boundary condition: One of

$$\beta_2 = \beta; \alpha_1 = \alpha \quad (\text{Dirichlet})$$

$$\beta_2 = (\sqrt{2})\beta; \alpha_1 = \alpha \quad (\text{Neumann centered})$$

$$\beta_2 = \beta; \alpha_1 = \alpha - \beta \quad (\text{Neumann noncentered})$$

$$\beta_2 = \psi \neq 0; \alpha_1 = \rho, \\ \rho \text{ arbitrary} \quad (\text{mixed}).$$

Bottom boundary condition: One of

$$\beta_N = \beta; \alpha_N = \alpha \quad (\text{Dirichlet})$$

$$\beta_N = (\sqrt{2})\beta; \alpha_N = \alpha \quad (\text{Neumann centered})$$

$$\beta_N = \beta; \alpha_N = \alpha - \beta \quad (\text{Neumann noncentered})$$

$$\beta_N = \chi \neq 0; \alpha_N = \eta, \\ \eta \text{ arbitrary} \quad (\text{mixed}).$$

NTYPE = 3: General separable; periodic boundary conditions.

$$\beta_i \neq 0, 1 \leq i \leq N;$$

$$\alpha_i \text{ arbitrary, } 1 \leq i \leq N.$$

NTYPE = 4: Constant coefficient; periodic boundary conditions.

$$\beta_i = \beta \neq 0, 1 \leq i \leq N;$$

$$\alpha_i = \alpha, \alpha \text{ arbitrary, } 1 \leq i \leq N.$$

If $N = 3$, GMA may respecify NTYPE.

TN1 is a real array of length $N + 1$, with $TN1(I) = \alpha_I, 1 \leq I \leq N$.

TN2 is a real array of length $N + 1$, with $TN2(I) = \beta_I, 1 \leq I \leq N$.

Input values of $TN2(N + 1)$, and $TN2(1)$ if $NTYPE \leq 2$, may be overwritten by GMA.

M is the dimension of the matrix TM; $M > 2$.

MTYPE is an integer describing the matrix TM.

MTYPE = 1: General separable; Dirichlet, Neumann, or mixed boundary conditions.

MTYPE = 2: Constant coefficient; Dirichlet, Neumann, or mixed boundary conditions.

MTYPE = 3: General separable; periodic boundary conditions.

MTYPE = 4: Constant coefficient; periodic boundary conditions.

Restrictions on γ_i and σ_i are analogous to those for α_i and β_i , respectively, for each given type. If $M \leq 5$, GMA may respecify MTYPE.

TM1 is a real array of length $M + 1$, with $TM1(I) = \gamma_I, 1 \leq I \leq M$.

TM2 is a real array of length $M + 1$, with $TM2(I) = \sigma_I, 1 \leq I \leq M$.

IBDIM is the row dimension of the array B, as it appears in the calling program.

B is a two-dimensional array with at least M rows and N columns. On input, B contains the right-hand side, with the vector b_I residing in the I-th column; $B(J, I) = (b_I)_J, 1 \leq J \leq M, 1 \leq I \leq N$. On output B contains the solution, with the vector x_I residing in the I-th column; $B(J, I) = (x_I)_J, 1 \leq J \leq M, 1 \leq I \leq N$.

K is the marching parameter. If $K < 2$, then K assumes the default value $K = 2$.

R is a real array of length $N \cdot P + 3 \cdot N + 2$, where $\log_2(N/K) \leq P$, P an integer, $K \geq 2$. The array R contains the output of a call to GMA with NPC = 0 (pre-processing calculations).

A is a real array of length $4 \cdot Q$, where $Q = \max(N + 1, M + 1)$. This array is used as a scratchpad array by GMA.

Subroutine GMA has one labeled common block, /MACHEP/, containing one

variable, TOL. TOL is a machine-dependent constant, which is equal to the machine epsilon. It is initialized in GMA in the first executable statement.

When the matrix A of eq. (1) is positive or negative *semidefinite*, with zero an eigenvalue of multiplicity one, then subroutine GMAS should be used in place of GMA to compute a least-squares solution. The usage of GMAS is nearly identical to that of GMA, though the two differ internally in several respects.

First, in the preprocessing phase ($NPC = 0$), the eigenvector v corresponding to the zero eigenvalue is determined by computing appropriate eigenvectors of the matrices TN and TM . The particular root S , for which $TM - SI$ is singular, is perturbed to guard against possible divide checks.

Second, in the backsolution phase ($NPC = 1$), the right-hand side b is projected into the orthogonal complement of the vector v . The right-hand sides for the reduced equations after STEP1 and STEP2 are also projected into their appropriate orthogonal complements, since the original orthogonalities may be diminished by roundoff growth during these steps of the backsolution. Finally, the solution x is projected into the orthogonal complement of v . This completes the least-squares solution.

Subroutine GMAS calls internally the 15 subroutines employed by GMA. Additionally, GMAS calls subroutines TINVIT, PINVIT, SVALUE, and PROJ as required. TINVIT and PINVIT are used to compute eigenvectors of symmetric tridiagonal matrices and symmetric matrices with periodic boundary conditions, respectively, using the method of inverse iteration [11]. TINVIT is a modified version of EISPACK subroutine of the same name [7, 12]. Subroutine SVALUE computes the perturbation of the root S . PROJ computes the projections described above.

Subroutine GMAS is entered using

```
CALL GMAS(NPC, N, NTYPE, TN1, TN2, M, MTYPE, TM1, TM2, IBDIM, B, VN, VM,
K, R, A).
```

All parameters except VN, VM, and A are identical to the corresponding parameters in the calling sequence for GMA.

VN is a real array of length $N + 1$. On output from a call to GMAS with $NPC = 0$,

VN contains the eigenvector of TN associated with the zero eigenvalue.

VM is a real array of length $M + 1$. On output from a call to GMAS with $NPC = 0$,

VM contains the eigenvector of TM associated with the zero eigenvalue.

A is a real array of length $5 \cdot Q$, where $Q = \max(N + 1, M + 1)$. This array is used as a scratchpad array by GMAS.

2. TESTS

The operation count for subroutine GMA (GMAS) is $O(N \cdot M \cdot NE) = O(N \cdot M \cdot \log_2(N/K))$, with the constant depending on the problem specification and the boundary conditions. (If $NE = 0$ this becomes $O(N \cdot M)$). Constant coefficient problems are important special cases of general separable problems in which the operation count can be reduced, often dramatically, by taking advantage of the additional restrictions on the coefficients. In Table I we have assembled execution times (in milliseconds) for GMA on an IBM 370-195 computer for some benchmark values of $N = M$ and K . Both constant coefficient and general separable problems were solved to illustrate the savings which can accrue in the special cases. The programs were compiled using the Fortran-H compiler, option = 2. Overall execution time is given as a sum of preprocessing time ($NPC = 0$) and backsolution time ($NPC = 1$).

The constant coefficient problems were all Helmholtz equations in the unit

square, with the Helmholtz constant a pseudorandom number between zero and one. Dirichlet boundary conditions were imposed on all boundaries. For the general separable problems, the β_i and σ_i , $1 \leq i \leq N + 1$, were all pseudorandom numbers between one and two. Then $\alpha_i = \beta_i + \beta_{i+1} + \rho_i$ and $\gamma_i = \sigma_i + \sigma_{i+1} + \eta_i$, $1 \leq i \leq N$, were calculated, with ρ_i and η_i pseudorandom numbers between zero and N^{-2} ; this insured that the resulting matrices were diagonally dominant, hence positive definite.

The savings which result in the constant coefficient case depend to some extent

Table I. Execution Times for Subroutine GMA on the IBM 370-195

(P = preprocessing time (NPC = 0) in milliseconds, B = backsolution time (NPC = 1) in milliseconds, T = total execution time in milliseconds. NTYPE = MTYPE = 1 in general separable problems, NTYPE = MTYPE = 2 in constant coefficient problems, N = M in all problems.)

	Constant coefficient problems			General separable problems			
	N = 31	N = 63	N = 127	N = 31	N = 63	N = 127	
K = 2	P	3	7	17	22	73	245
	B	24	103	449	44	215	1033
	T	27	110	446	66	288	1278
K = 4	P	3	6	14	22	72	243
	B	22	99	444	37	191	949
	T	25	105	458	59	263	1192
K = 8	P	2	5	12	19	68	234
	B	18	86	401	27	153	806
	T	20	91	413	46	221	1040
K = 16	P	1	4	9	15	59	217
	B	14	72	350	17	112	633
	T	15	76	359	32	171	850
K = 32	P	1	2	7	9	48	194
	B	8	50	287	9	70	457
	T	9	52	294	18	118	651

on the boundary conditions, the Dirichlet problem illustrated being the optimal case. For both constant coefficient and general separable problems, the execution times increase if periodic boundary conditions are imposed.

In Table II we present some results illustrating the numerical stability of the generalized marching algorithm as a function of $N = M$ and K for each problem type. An exact solution x_e of pseudorandom numbers on $[-1, 1]$ was generated and substituted into the difference equation in order to generate a right-hand side. The linear system was then solved, and the computed solution x_c was compared with the exact solution. The number of correct digits was computed using $\text{digits} = -\log_{10} ((e^T e / x_e^T x_e)^{1/2})$, $e = x_c - x_e$. REAL*8 arithmetic was used (approximately 16 decimal digits).

For $K = 2, 4$, the error primarily reflects the condition number of the matrix A . When $K = 8, 16, 32$, however, an exponential "marching" term becomes the dominant term in the error [1].

Note that both execution time and numerical stability decrease with increasing K ; thus one must seek to strike a balance between rapid execution times and accuracy requirements.

Table II. Numerical Stability of Subroutine GMA on the IBM 370-195

(NTYPE = MTYPE = 1 in general separable problems, NTYPE = MTYPE = 2 in constant coefficient problems, N = M in all problems.)

	constant coefficient problems			general separable problems		
	N = 31	N = 63	N = 127	N = 31	N = 63	N = 127
K = 2	14.8	14.3	13.9	13.3	12.5	12.6
K = 4	14.2	14.0	13.9	13.3	12.5	12.6
K = 8	11.1	11.1	11.0	11.1	11.0	11.6
K = 16	4.9	4.9	4.9	4.6	4.1	4.0
K = 32	0.0	0.0	0.0	0.0	0.0	0.0

Subroutine KPICK is designed to aid the user in making appropriate choices of the marching parameter K and to provide other information useful in determining if a linear system is appropriate for GMA or GMAS.

KPICK can be used in two ways: first, to find an optimal value of K for a user-specified accuracy demand, or, second, to estimate the accuracy for a user-selected value of K.

KPICK is based on a heuristic which supposes that subroutine GMA (or GMAS) yields solutions satisfying error estimates of the form $(\text{COND} + \text{EMARCH}) * \text{TOL}$, where COND is the condition number of the linear system EMARCH is the marching error, and TOL is the machine epsilon. In [1], error bounds of this form are proved for constant coefficient Dirichlet problems when $N = K^{2^L} - 1$ for a non-negative integer L.

Subroutine KPICK is entered using

```
CALL KPICK(NPC, N, NTYPE, TN1, TN2, M, MTYPE, TM1, TM2, DEMAND, K, COND,
EMARCH, DIGITS, IFLAG)
```

The parameter list is:

NPC is an integer. If NPC = 0, KPICK determines K such that the solutions computed using subroutine GMA or GMAS will have approximately DEMAND significant digits. If NPC = 1, subroutine will test the input value of K.

N, NTYPE, TN1, TN2, M, MTYPE, TM1, TM2 are identical to the corresponding parameters in the calling sequence for subroutine GMA.

DEMAND is the real number, stating the number of significant digits desired in the computed solution. It must be specified if NPC = 0.

K is an integer. If NPC = 0, on output K is equal to the marching parameter determined by KPICK. If NPC = 1, the user must specify the value of K to be tested as an input parameter.

COND is a real number, normally returning the condition number of the linear system.

EMARCH is a real number, normally returning an estimate of the marching error for the output value of K.

DIGITS is a real number, normally returning a value of $\max(0.0, -\log((\text{COND} + \text{EMARCH}) * \text{TOL}))$. This is an estimate of the number of significant digits, in the 2-norm, one may expect in the computed solution.

IFLAG is an integer, describing error returns.

IFLAG = 0: Normal return.

IFLAG = 1: KPICK failed to successfully compute COND. The linear system is not positive or negative definite. If the linear system is positive or negative semidefinite with a zero eigenvalue of multiplicity one, the condition is computed

relative to the orthogonal complement of the eigenvector associated with the zero eigenvalue. Otherwise the default values $COND = EMARCH = DIGITS = 0.0$ are returned.

IFLAG = 2: The marching error may not satisfy the assumptions underlying the algorithm used to compute EMARCH. (The marching error appears to grow more slowly than $(2.5)^K$.) If NPC = 0, KPICK returns estimates for $K = 2$; if NPC = 1, KPICK returns estimates for the input value of K.

IFLAG = 3: Conditions 1 and 2 exist.

IFLAG = 4: The linear system cannot be solved to DEMAND significant digits for any $K \geq 2$. KPICK returns estimates for $K = 2$. This return can occur only if NPC = 0.

IFLAG = 5: Conditions 1 and 4 exist.

IFLAG = 6: Conditions 2 and 4 exist.

IFLAG = 7: Conditions 1, 2, and 4 exist.

IFLAG = 8: The input values of N, NTYPE, TN1, and/or TN2 are incorrectly specified.

IFLAG = 9: The input values of M, MTYPE, TM1, and/or TM2 are incorrectly specified.

IFLAG = 10: Conditions 8 and 9 exist.

If IFLAG = 0, 2, 4, 6, then the problem is appropriate for subroutine GMA with the qualifications noted above. If IFLAG = 1, 3, 5, 7, with other than default values for COND, EMARCH, and DIGITS, then the problem is appropriate for GMAS.

Subroutine KPICK calls subroutines TCHECK, EIGEN, TRIEIG, PEREIG, and ERROR as required. TCHECK is used to determine if the matrices TN and TM have been correctly specified. Then COND is determined by computing the extremal eigenvalues of TN and TM; this is done in EIGEN. For some matrix types, the eigenvalues are well known; otherwise TRIEIG or PEREIG are used. These routines compute eigenvalues of symmetric tridiagonal and periodic matrices, respectively, using bisection and the Sturm sequence property [4, 11].

For a fixed value of K, EMARCH is defined to be the largest 2-norm over the set of marching polynomials arising in STEP1 or STEP4 of the generalized marching algorithm [1, 2]. For a fixed value of K, subroutine ERROR is used to compute the marching error; the three-term recurrence relation for generating the marching polynomials is employed [1, 2]. If NPC = 1 computation of EMARCH is straightforward; if NPC = 0, the method of bisection is used to determine the optimal value of K, and hence EMARCH.

To test the effectiveness of KPICK, 3000 random problems (1000 nonsingular, 1000 singular, and 1000 where singularity was randomly determined) were solved on an IBM 370-195 using REAL*8 arithmetic. All required input values for KPICK, GMA, and GMAS were generated using random number generators of the type used in testing the EISPACK codes [7].

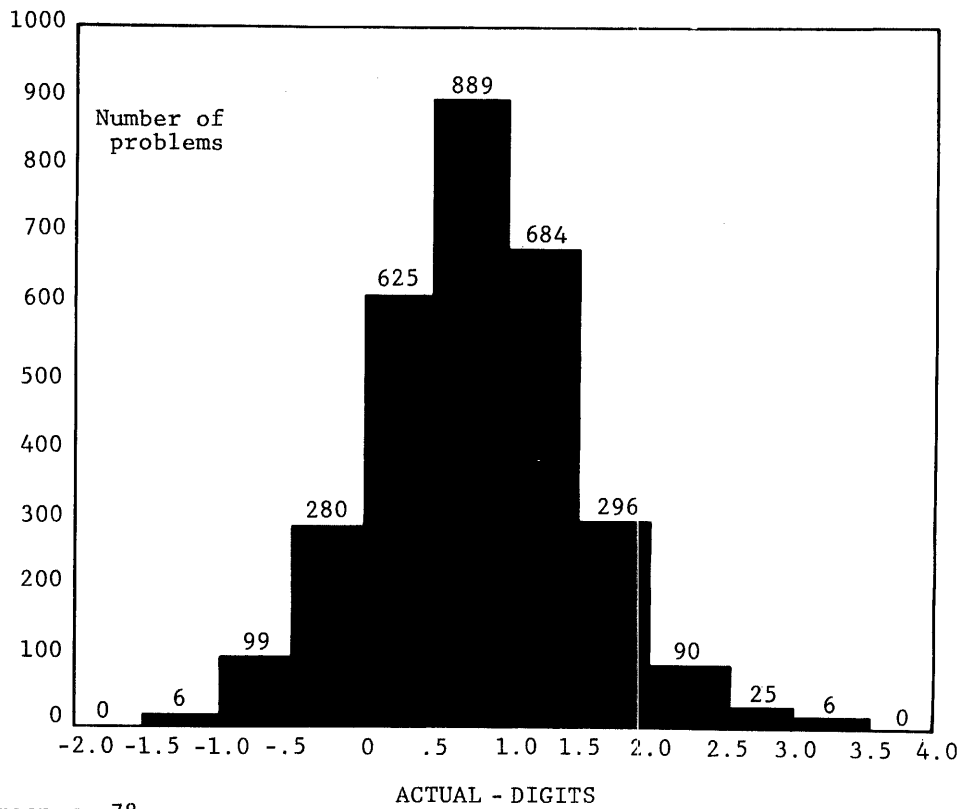
The values of N and M were random integers between 3 and 63; NTYPE and MTYPE were random integers between 1 and 4 (for type-2 matrices, the boundary condition combination was determined randomly). The matrix entries themselves were generated using random real numbers on $[10^{-3}, 1.0]$. Positive or negative definiteness (semidefiniteness) was determined randomly.

The parameter DEMAND was a random real number on $[1.0, -\log(TOL) - 1.0]$. Exact solutions were generated using random numbers on $[-1.0, 1.0]$ and were substituted into the difference equations to generate right-hand sides.

The random problems were analyzed by KPICK, with NPC = 0; the values of K produced in this fashion ranged between 2 and 29. The problems were then solved, using GMA or GMAS, and the computed solutions compared with the exact solutions. The number of correct digits, in the 2-norm, of the computed solution ACTUAL was determined and compared with the estimate of DIGITS provided by KPICK. The results are summarized in Table III, where we have

Table III. Effectiveness of Subroutine KPICK in Predicting the Error for
3000 Random Problems

(ACTUAL = actual number of correct digits in the computed solution in the 2-norm using GMA or GMAS, DIGITS = number of correct digits predicted by KPICK.)



mean = .78

variance = .48

graphed the value of ACTUAL-DIGITS versus the number of problems. The results indicate the effectiveness of KPICK in predicting the accuracy of computed solutions using GMA and GMAS. The results for singular and nonsingular problems considered separately are essentially the same as those given in Table III.

Other Fortran programs for solving eq. (1) or certain special cases using fast direct methods can be found in [3, 8]; theoretical discussions of the algorithms are given in [5, 6, 9, 10].

ACKNOWLEDGMENT

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ALGORITHM

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SUBROUTINE GMA(NPC,N,NTYPE,TN1,TN2,M,MTYPE,          GMA00010
1      TM1,TM2,IBDIM,B,K,R,A)                       GMA00020
C      SUBROUTINE GMA IS AN IMPLEMENTATION OF THE GENERALIZED GMA00030
C      MARCHING ALGORITHM, WHICH MAY BE USED FOR SOLVING LINEAR GMA00040
C      SYSTEMS ARISING FROM 5 POINT DISCRETIZATIONS OF SEPARABLE GMA00050
C      OR CONSTANT COEFFICIENT ELLIPTIC BOUNDARY VALUE PROBLEMS GMA00060
      ON RECTANGULAR DOMAINS. A DIRICHLET, NEUMANN, OR MIXED GMA00070
      BOUNDARY CONDITION MAY BE INDEPENDENTLY SPECIFIED ON EACH GMA00080
      SIDE OF THE RECTANGLE, OR PERIODIC BOUNDARY CONDITIONS MAY GMA00090
      BE SPECIFIED ON OPPOSING SIDES. GMA00100
      THE LINEAR SYSTEM HAS THE FORM: GMA00110
      GMA00120
      ( TM1(I) + TN1(J) ) * X(I,J) GMA00130
      - TM2(I+1) * X(I+1,J) - TM2(I) * X(I-1,J) GMA00140
      - TN2(J+1) * X(I,J+1) - TN2(J) * X(I,J-1) = B(I,J) GMA00150
      GMA00160
      FOR I = 1,2,...,M, AND J = 1,2,...,N, WHERE TM2(M+1) IS GMA00170
      INTERPRETED AS TM2(1), TN2(N+1) AS TN2(1), X(0,J) AS X(M,J), GMA00180
      X(M+1,J) AS X(1,J), X(I,0) AS X(I,N), AND X(I,N+1) AS X(I,1). GMA00190
      GMA00200
      THE PARAMETER LIST: GMA00210
      GMA00220
      NPC IS AN INTEGER. IF NPC = 0, GMA DOES NECESSARY GMA00230
      PREPROCESSING CALCULATIONS. IF NPC = 1, GMA SOLVES THE GMA00240
      LINEAR SYSTEM. GMA00250
      N IS AN INTEGER, AS DEFINED ABOVE. N MUST BE GREATER THAN 2. GMA00260
      NTYPE IS AN INTEGER DESCRIBING THE ARRAYS TN1 AND TN2. GMA00270
      NTYPE = 1: GENERAL SEPARABLE - DIRICHLET, NEUMANN, OR GMA00280
      MIXED BOUNDARY CONDITIONS. GMA00290
      TN1(I) IS ARBITRARY, I = 1,2,...,N. GMA00300
      TN2(1) = 0.0; TN2(I) IS ARBITRARY, NONZERO, I = 2,3,...,N. GMA00310
      NTYPE = 2: CONSTANT COEFFICIENT - DIRICHLET, NEUMANN, GMA00320
      OR MIXED BOUNDARY CONDITIONS. GMA00330
      TN1(I) = ALPHA, ALPHA AN ARBITRARY CONSTANT, I = 2,3,...,N-1. GMA00340
      TN2(1) = 0.0; TN2(I) = BETA, BETA AN ARBITRARY NONZERO GMA00350
      CONSTANT, I = 3,...,N-1. GMA00360
      TOP BOUNDARY CONDITION : ONE OF GMA00370
      TN1(1) = ALPHA; TN2(2) = BETA; (DIRICHLET) GMA00380
      TN1(1) = ALPHA; TN2(2) = SQRT(2) * BETA; (NEUMANN-CENTERED) GMA00390
      TN1(1) = ALPHA - BETA; TN2(2) = BETA; (NEUMANN-STAGGERED) GMA00400
      TN1(1) = RHO, RHO ARBITRARY; TN2(2) = ZETA, ZETA ARBITRARY, GMA00410
      NONZERO; (MIXED). GMA00420
      BOTTOM BOUNDARY CONDITION : ONE OF GMA00430
      TN1(N) = ALPHA; TN2(N) = BETA; (DIRICHLET) GMA00440
      TN1(N) = ALPHA; TN2(N) = SQRT(2) * BETA; (NEUMANN-CENTERED) GMA00450
      TN1(N) = ALPHA - BETA; TN2(N) = BETA; (NEUMANN-STAGGERED) GMA00460
      TN1(N) = CHI, CHI ARBITRARY; TN2(N) = ETA, ETA ARBITRARY, GMA00470
      NONZERO; (MIXED). GMA00480
      NTYPE = 3: GENERAL SEPARABLE - PERIODIC BOUNDARY CONDITIONS. GMA00490
      TN1(I) IS ARBITRARY, I = 1,2,...,N. GMA00500
      TN2(I) IS ARBITRARY, NONZERO, I = 1,2,...,N. GMA00510
      NTYPE = 4: CONSTANT COEFFICIENT - PERIODIC BOUNDARY GMA00520
      CONDITIONS. GMA00530
      TN1(I) = ALPHA, ALPHA ARBITRARY, I = 1,2,...,N. GMA00540
      TN2(I) = BETA, BETA ARBITRARY, NONZERO, I = 1,2,...,N. GMA00550
      IF N = 3, GMA MAY RESPECIFY NTYPE. GMA00560
      TN1 IS AN ARRAY OF LENGTH N+1, DEFINED ABOVE. GMA00570
      TN2 IS AN ARRAY OF LENGTH N+1, DEFINED ABOVE. GMA00580
      M IS AN INTEGER, AS DEFINED ABOVE. M MUST BE GREATER THAN 2. GMA00590

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MTYPE IS AN INTEGER DESCRIBING THE ARRAYS TM1 AND TM2.          GMA00600
  MTYPE = 1: GENERAL SEPARABLE - DIRICHLET, NEUMANN, OR        GMA00610
  MIXED BOUNDARY CONDITIONS.                                  GMA00620
  MTYPE = 2: CONSTANT COEFFICIENT - DIRICHLET, NEUMANN,      GMA00630
  OR MIXED BOUNDARY CONDITIONS.                              GMA00640
  MTYPE = 3: GENERAL SEPARABLE - PERIODIC BOUNDARY CONDITIONS. GMA00650
  MTYPE = 4: CONSTANT COEFFICIENT - PERIODIC BOUNDARY        GMA00660
  CONDITIONS.                                                GMA00670
RESTRICTIONS ON TM1 AND TM2 ARE ANALOGOUS TO THOSE ON        GMA00680
  TN1 AND TN2, RESPECTIVELY, FOR EACH GIVEN TYPE.           GMA00690
  IF M IS LESS THAN 6, GMA MAY RESPECIFY MTYPE.              GMA00700
TM1 IS AN ARRAY OF LENGTH M+1, DEFINED ABOVE.                GMA00710
C   TM2 IS AN ARRAY OF LENGTH M+1, DEFINED ABOVE.            GMA00720
C   IBDIM IS THE ROW DIMENSION OF THE ARRAY B, AS IT APPEARS GMA00730
C   IN THE CALLING PROGRAM.                                  GMA00740
C   B IS A TWO DIMENSIONAL ARRAY WITH AT LEAST M ROWS AND N  GMA00750
C   COLUMNS. ON INPUT, B CONTAINS THE RIGHT HAND SIDE B(I,J) GMA00760
C   AS DEFINED ABOVE. ON OUTPUT, FROM A CALL TO GMA WITH     GMA00770
C   NPC = 1, B CONTAINS THE SOLUTION X(I,J), AS DEFINED ABOVE. GMA00780
C   K IS THE MARCHING PARAMETER. MARCHING OCCURS IN THE     GMA00790
C   N - DIRECTION. ON INPUT, IF K IS LESS THAN 2, IT        GMA00800
C   ASSUMES THE DEFAULT VALUE K = 2.                          GMA00810
C   R IS AN ARRAY OF LENGTH N * P + 3 * N + 2, WHERE P IS AN GMA00820
C   INTEGER GREATER THAN OR EQUAL TO LOG2( N / K ), AND K IS GMA00830
C   GREATER THAN OR EQUAL TO 2. R CONTAINS THE OUTPUT FROM A GMA00840
C   CALL TO GMA WITH NPC = 0.                                  GMA00850
C   A IS AN ARRAY OF LENGTH 4 * Q, WHERE Q = MAX( N+1, M+1 ). GMA00860
C   A IS USED AS A SCRATCHPAD ARRAY BY GMA.                   GMA00870
C                                                             GMA00880
C   SUBROUTINE GMA HAS ONE LABELED COMMON BLOCK, /MACHEP/,    GMA00890
C   WITH ONE VARIABLE, TOL. TOL IS A MACHINE DEPENDENT CONSTANT, GMA00900
C   EQUAL TO THE MACHINE EPSILON. IT IS INITIALIZED IN GMA IN THE GMA00910
C   FIRST EXECUTABLE STATEMENT.                               GMA00920
C                                                             GMA00930
C   A CALL TO GMA WITH NPC = 0 MUST PRECEDE A CALL TO GMA    GMA00940
C   WITH NPC = 1. IF THE SAME LINEAR SYSTEM IS TO BE SOLVED  GMA00950
C   WITH SEVERAL RIGHT HAND SIDES, A SINGLE CALL TO GMA WITH GMA00960
C   NPC = 0 WILL SUFFICE, PROVIDED THAT THE CONTENTS OF THE  GMA00970
C   ARRAY R ARE SAVED.                                        GMA00980
C   SUBROUTINE GMA MAY FAIL TO GIVE CORRECT ANSWERS IF       GMA00990
C   THE LINEAR SYSTEM IS NOT POSITIVE OR NEGATIVE DEFINITE.  GMA01000
C   IF N, NTYPE, M, MTYPE, OR IBDIM ARE OUT OF RANGE, GMA RETURNS GMA01010
C   WITHOUT ATTEMPTING ANY COMPUTATIONS.                      GMA01020
C                                                             GMA01030
C   ADDRESS INQUIRIES TO:                                     GMA01040
C   RANDOLPH E. BANK                                         GMA01050
C   DEPARTMENT OF MATHEMATICS                                 GMA01060
C   THE UNIVERSITY OF CHICAGO                                 GMA01070
C   CHICAGO ILLINOIS 60637.                                  GMA01080
C                                                             GMA01090
C   VERSION DATE: MARCH 1, 1977.                              GMA01100
C   DIMENSION TN1(1),TN2(1),TM1(1),TM2(1),A(1),B(1),R(1)    GMA01110
C   COMMON /MACHEP/TOL                                       GMA01120
C   TOL=2.0E0**(-23)                                          GMA01130
C   CHECK N, NTYPE, M, MTYPE, AND IBDIM.                     GMA01140
C   IF((NTYPE.GT.4).OR.(NTYPE.LT.1)) RETURN                  GMA01150
C   IF((MTYPE.GT.4).OR.(MTYPE.LT.1)) RETURN                  GMA01160
C   IF((N.LT.2).OR.(M.LT.2)) RETURN                          GMA01170
C   IF(IBDIM.LT.M) RETURN                                     GMA01180
C   RESPECIFY NTYPE AND/OR MTYPE IF NECESSARY.                GMA01190
C   IF(NTYPE.EQ.2.AND.N.LE.3) NTYPE=1                         GMA01200
C   IF(MTYPE.EQ.2.AND.M.LE.3) MTYPE=1                         GMA01210
C   IF(NTYPE.EQ.4.AND.N.LE.3) NTYPE=3                         GMA01220
C   IF(MTYPE.EQ.4.AND.M.LE.5) MTYPE=3                         GMA01230
C   IF(N.LE.2) NTYPE=1                                         GMA01240
C   IF(M.LE.2) MTYPE=1                                         GMA01250
C   IF(NPC.NE.0) GO TO 15                                      GMA01260
C   THE PREPROCESSING CALCULATIONS.                            GMA01270
C   CALL PARTN(N,NTYPE,K,NE,ND,R(3))                          GMA01280
C   EPS IS ADDED TO ALL INTEGERS STORED IN THE REAL ARRAY R. GMA01290
C   THIS PREVENTS UNFORTUNATE ROUNDING ERRORS FORM OCCURRING WHEN GMA01300
C   THEY ARE CONVERTED BACK INTO INTEGERS.                    GMA01310
C   EPS=0.25E0                                                GMA01320
C   R(1)=FLOAT(NE)+EPS                                         GMA01330
C   R(2)=FLOAT(ND)+EPS                                         GMA01340
C   IF(NTYPE.EQ.1.OR.NTYPE.EQ.3) CALL ROOTSG(N,NTYPE,        GMA01350
C   TN1,TN2,NE,ND,R(3),R(N+3),A)                             GMA01360
C   IF(NTYPE.EQ.2.OR.NTYPE.EQ.4) CALL ROOTSC(N,NTYPE,        GMA01370

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1      TN1,TN2,NE,ND,R(3),R(N+3),A)
      J1=N+2
      J2=J1+N+1
      CALL SORT(N,NTYPE,NE,ND,R(3),R(N+3),A(1),A(J1),A(J2))
      RETURN
C     THE BACKSOLUTION CALCULATIONS.
15    TN2(N+1)=1.0E0
      IF(NTYPE.LT.3) TN2(1)=1.0E0
      NE=R(1)
      ND=R(2)
      CALL STEP1(N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2,
1      IBDIM,B,ND,R(3),R(N+3),A)
      CALL STEP2(N,NTYPE,TN2,M,MTYPE,TM1,TM2,
1      IBDIM,B,NE,ND,R(3),R(N+3),A)
      CALL STEP3(N,NTYPE,TN2,M,MTYPE,TM1,TM2,
1      IBDIM,B,NE,ND,R(3),R(N+3),A)
      CALL STEP4(TN1,TN2,M,MTYPE,TM1,TM2,IBDIM,B,ND,R(3),R(N+3),A)
      RETURN
      END
C
C
      SUBROUTINE PARTN(N,NTYPE,K,NE,ND,ID)
C     SUBROUTINE PARTN PARTITIONS THE GRID, IN THE N - DIRECTION,
C     INTO ND+1 BLOCKS, EACH WITH K OR FEWER LINES. POINTERS TO THE
C     SEPARATING LINES ARE STORED IN THE ARRAY ID. THE VALUE
C     OF THE INTEGER NE, ROUGHLY EQUAL TO LOG2(ND), IS COMPUTED.
C     EPS IS ADDED TO ALL INTEGERS STORED IN THE REAL ARRAY ID.
C     THIS PREVENTS UNFORTUNATE ROUNDING ERRORS FORM OCCURRING WHEN
C     THEY ARE CONVERTED BACK INTO INTEGERS.
C     VERSION DATE: FEBRUARY 1, 1977.
      REAL ID
      DIMENSION ID(1)
      EPS=.25E0
      NN=N
      IF(NTYPE.GE.3) NN=N-1
      ND=NN/MAX0(K,2)
      NE=0
      IF(2*ND.EQ.NN) ND=ND-1
      IF(ND.EQ.0) GO TO 10
      I1=NN/(ND+1)
      I2=NN-I1*(ND+1)
      I3=0
      DO 5 J=1,ND
          IF(I3.LT.I2) I3=I3+1
          J1=J*I1+I3
          ID(J+1)=FLOAT(J1)+EPS
          J2=2** (J-1)
          IF(J2.LE.ND.AND.ND.LE.(2*J2-1)) NE=J
      5      CONTINUE
10     ID(1)=EPS
      I1=ND+2
      DO 15 J=I1,N
15     ID(J)=FLOAT(NN+1)+EPS
      RETURN
      END
C
C
      SUBROUTINE ROOTSC(N,NTYPE,TN1,TN2,NE,ND,ID,R,A)
C     SUBROUTINE ROOTSC DOES THE NECESSARY EIGENVALUE
C     CALCULATIONS FOR CONSTANT COEFFICIENT PROBLEMS.
C     THE EIGENVALUES ARE STORED IN THE ARRAY R.
C     VERSION DATE: FEBRUARY 1, 1977.
      REAL ID
      DIMENSION TN1(1),TN2(1),ID(1),R(1),A(1)
      NP1=N+1
      NEP1=NE+1
      NDP1=ND+1
      PI=3.1415926535897932384626433832795E0
      T1=TN1(2)
      T2=TN2(3)
      T22=2.0E0*T2
      S=ABS(T22)
C     IN THE FOLLOWING SEGMENT OF CODE, SPECIAL EIGENVALUES FOR THE
C     CASE NTYPE = 4 ARE COMPUTED.
      IF(NTYPE.EQ.2) GO TO 130
      LINE=N*NEP1
      I2=LINE+N

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GMA01380
GMA01390
GMA01400
GMA01410
GMA01420
GMA01430
GMA01440
GMA01450
GMA01460
GMA01470
GMA01480
GMA01490
GMA01500
GMA01510
GMA01520
GMA01530
GMA01540
GMA01550
GMA01560
PRTN0010
PRTN0020
PRTN0030
PRTN0040
PRTN0050
PRTN0060
PRTN0070
PRTN0080
PRTN0090
PRTN0100
PRTN0110
PRTN0120
PRTN0130
PRTN0140
PRTN0150
PRTN0160
PRTN0170
PRTN0180
PRTN0190
PRTN0200
PRTN0210
PRTN0220
PRTN0230
PRTN0240
PRTN0250
PRTN0260
PRTN0270
PRTN0280
PRTN0290
PRTN0300
PRTN0310
PRTN0320
PRTN0330
PRTN0340
RTSC0010
RTSC0020
RTSC0030
RTSC0040
RTSC0050
RTSC0060
RTSC0070
RTSC0080
RTSC0090
RTSC0100
RTSC0110
RTSC0120
RTSC0130
RTSC0140
RTSC0150
RTSC0160
RTSC0170
RTSC0180
RTSC0190
RTSC0200

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R(I2)=T1-S
R(LINE+1)=T1+S
NM1=N-1
I2=2
Q=PI/FLOAT(N)
IF(T22.LT.0.0E0) GO TO 110
LINE=LINE+N
I2=-2
110 DO 115 I=2,NM1,2
    LINE=LINE+I2
    ARG=Q*FLOAT(I)
    R(LINE)=T1-T22*COS(ARG)
115 R(LINE+1)=R(LINE)
C IN THE FOLLOWING SEGMENT OF CODE, EIGENVALUES RELEVANT TO ALL
C CONSTANT COEFFICIENT PROBLEMS ARE CALCULATED.
130 DO 140 I=1,NEP1
    I2=2*(I-1)
    J2=N*(I-1)
    DO 140 LNINDEX=1,NDP1,I2
        LINE=ID(LNINDEX)
        ISPAN=LNINDEX+I2
        ISPAN=ID(ISPAN)
        ISPAN=ISPAN-LINE-1
        LINE=LINE+J2
        Q=PI/FLOAT(ISPAN+1)
        DO 135 J=1,ISPAN
            LINE=LINE+1
            ARG=Q*FLOAT(J)
135 R(LINE)=T1+S*COS(ARG)
140 R(LINE+1)=0.0E0
    CONTINUE
IF(NTYPE.EQ.4) RETURN
C IN THE FOLLOWING SEGMENT OF CODE, WE CHECK FOR NEUMANN OR MIXED
C BOUNDARY CONDITIONS, AND MODIFY SOME EIGENVALUES IN THE
C APPROPRIATE FASHION.
C ICOUNT KEEPS TRACK OF THE BOUNDARY CONDITION COMBINATION.
ICOUNT=0
IA1=1
IA2=0
S2=SQRT(2.0E0)*T2
DO 270 L=1,2
    I2=2*IA1+N*IA2
    T2P=TN2(I2)
    I2=I2-IA1
    T1P=TN1(I2)
    IF(T2P.EQ.T2.AND.T1P.EQ.T1) GO TO 265
    NN=ID(ND+2)
    NN=NN*IA2
    IF(T1P.EQ.T1.AND.T2P.EQ.S2) GO TO 220
    IF(T2P.EQ.T2.AND.T1P.EQ.(T1-T2)) GO TO 225
C THE LOOP FOR A MIXED BOUNDARY CONDITION
ICOUNT=ICOUNT+7
IF(NE.EQ.0) GO TO 265
DO 215 I=1,NE
    I2=2*(I-1)
    LINE=(ND/I2)*IA2*I2
    LINE=ID(LINE+1)
    ISPAN=ID(I2+1)
    ISPAN=ISPAN*IA1+NN-LINE-1
    LINE=LINE+N*(I-1)
    J2=N
    DO 205 J=1,ISPAN
        J2=J2+1
        A(J)=T1
205 A(J2)=-T2
        A(J2)=0.0E0
        A(1)=T1P
        A(NP1)=-T2P
        CALL QL(A(1),A(NP1),ISPAN,IERR)
        J2=ISPAN+1
        DO 210 J=1,ISPAN
            J2=J2-1
            LINE=LINE+1
210 R(LINE)=A(J2)
215 CONTINUE
GO TO 265
C THE LOOP FOR A NEUMANN BOUNDARY CONDITION
220 ICOUNT=ICOUNT+1

```

```

RTSC0210
RTSC0220
RTSC0230
RTSC0240
RTSC0250
RTSC0260
RTSC0270
RTSC0280
RTSC0290
RTSC0300
RTSC0310
RTSC0320
RTSC0330
RTSC0340
RTSC0350
RTSC0360
RTSC0370
RTSC0380
RTSC0390
RTSC0400
RTSC0410
RTSC0420
RTSC0430
RTSC0440
RTSC0450
RTSC0460
RTSC0470
RTSC0480
RTSC0490
RTSC0500
RTSC0510
RTSC0520
RTSC0530
RTSC0540
RTSC0550
RTSC0560
RTSC0570
RTSC0580
RTSC0590
RTSC0600
RTSC0610
RTSC0620
RTSC0630
RTSC0640
RTSC0650
RTSC0660
RTSC0670
RTSC0680
RTSC0690
RTSC0700
RTSC0710
RTSC0720
RTSC0730
RTSC0740
RTSC0750
RTSC0760
RTSC0770
RTSC0780
RTSC0790
RTSC0800
RTSC0810
RTSC0820
RTSC0830
RTSC0840
RTSC0850
RTSC0860
RTSC0870
RTSC0880
RTSC0890
RTSC0900
RTSC0910
RTSC0920
RTSC0930
RTSC0940
RTSC0950
RTSC0960
RTSC0970
RTSC0980

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```

INS=0
GO TO 230
225 ICOUNT=ICOUNT+3
INS=1
230 IF(NE.EQ.0) GO TO 265
J1=0
J2=2
IF(T22.LT.0.0E0) GO TO 240
J1=2
J2=-2
240 DO 250 I=1,NE
I2=2**(I-1)
LINE=(ND/I2)*IA2*I2
LINE=ID(LINE+1)
ISPAN=ID(I2+1)
ISPAN=ISPAN*IA1+NN-LINE-1
LINE=LINE+N*(I-1)
Q=PI/FLOAT(2*ISPAN+INS)
J=J1*(ISPAN+1)-1
DO 245 JJ=1,ISPAN
J=J+J2
LINE=LINE+1
ARG=Q*FLOAT(J)
245 R(LINE)=T1-T22*COS(ARG)
250 CONTINUE
265 IA1=0
270 IA2=1
IF(ICOUNT.EQ.0) RETURN
LINE=N*NE
IF(ICOUNT.GE.7) GO TO 400
GO TO (425,430,435,440,400,445),ICOUNT
C IN THE FOLLOWING SEGMENT, EIGENVALUES FOR THE CASE OF ONE
C OR MORE MIXED BOUNDARY CONDITIONS ARE CALCULATED.
400 J2=N
DO 405 J=1,N
J2=J2+1
A(J)=T1
405 A(J2)=-T2
A(1)=TN1(1)
A(N)=TN1(N)
A(NP1)=-TN2(2)
A(J2-1)=-TN2(N)
A(J2)=0.0E0
CALL QL(A(1),A(NP1),N,IERR)
J2=NP1
DO 410 J=1,N
LINE=LINE+1
J2=J2-1
410 R(LINE)=A(J2)
RETURN
C IN THE FOLLOWING SEGMENT, EIGENVALUES FOR THE CASE OF ONE
C NEUMANN AND ONE DIRICHLET BOUNDARY CONDITION ARE CALCULATED.
425 Q=PI/FLOAT(2*N)
T22=-S
J=-1
J2=2
GO TO 450
C IN THE FOLLOWING SEGMENT, EIGENVALUES FOR THE CASE OF TWO
C NEUMANN BOUNDARY CONDITIONS ARE CALCULATED.
430 Q=PI/FLOAT(N-1)
T22=-S
J=-1
J2=1
GO TO 450
C IN THE FOLLOWING SEGMENT, EIGENVALUES FOR THE CASE OF ONE
C STAGGERED AND ONE DIRICHLET BOUNDARY CONDITION ARE CALCULATED.
435 Q=PI/FLOAT(2*N+1)
J=-1
J2=2
IF(T22.LT.0.0E0) GO TO 450
J=2*N+1
J2=-2
GO TO 450
C IN THE FOLLOWING SEGMENT, EIGENVALUES FOR THE CASE OF ONE
C NEUMANN AND ONE STAGGERED BOUNDARY CONDITION ARE CALCULATED.
440 Q=PI/FLOAT(2*N-1)

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RTSC0990
RTSC1000
RTSC1010
RTSC1020
RTSC1030
RTSC1040
RTSC1050
RTSC1060
RTSC1070
RTSC1080
RTSC1090
RTSC1100
RTSC1110
RTSC1120
RTSC1130
RTSC1140
RTSC1150
RTSC1160
RTSC1170
RTSC1180
RTSC1190
RTSC1200
RTSC1210
RTSC1220
RTSC1230
RTSC1240
RTSC1250
RTSC1260
RTSC1270
RTSC1280
RTSC1290
RTSC1300
RTSC1310
RTSC1320
RTSC1330
RTSC1340
RTSC1350
RTSC1360
RTSC1370
RTSC1380
RTSC1390
RTSC1400
RTSC1410
RTSC1420
RTSC1430
RTSC1440
RTSC1450
RTSC1460
RTSC1470
RTSC1480
RTSC1490
RTSC1500
RTSC1510
RTSC1520
RTSC1530
RTSC1540
RTSC1550
RTSC1560
RTSC1570
RTSC1580
RTSC1590
RTSC1600
RTSC1610
RTSC1620
RTSC1630
RTSC1640
RTSC1650
RTSC1660
RTSC1670
RTSC1680
RTSC1690
RTSC1700
RTSC1710
RTSC1720
RTSC1730
RTSC1740

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J=-2
J2=2
IF(T22.LT.0.0E0) GO TO 450
J=2*N
J2=-2
GO TO 450
C IN THE FOLLOWING SEGMENT, EIGENVALUES FOR THE CASE OF TWO
C STAGGERED BOUNDARY CONDITIONS ARE CALCULATED.
445 Q=PI/FLOAT(N)
J=-1
J2=1
IF(T22.LT.0.0E0) GO TO 450
J=N
J2=-1
450 DO 455 JJ=1,N
J=J+J2
LINE=LINE+1
ARG=Q*FLOAT(J)
455 R(LINE)=T1-T22*COS(ARG)
RETURN
END

C
C SUBROUTINE ROOTSG(N,NTYPE,TN1,TN2,NE,ND,ID,R,A)
C SUBROUTINE ROOTSG DOES THE NECESSARY EIGENVALUE
C CALCULATIONS FOR GENERAL SEPARABLE PROBLEMS.
C THE EIGENVALUES ARE STORED ON THE ARRAY R.
C VERSION DATE: FEBRUARY 1, 1977.
REAL ID
DIMENSION TN1(1),TN2(1),ID(1),R(1),A(1)
NP1=N+1
NEP1=NE+1
NDP1=ND+1
C IN THE FOLLOWING SEGMENT OF CODE, SPECIAL EIGENVALUES FOR THE
C CASE NTYPE = 3 ARE COMPUTED.
IF(NTYPE.EQ.1) GO TO 50
N2=(NP1)/2
NP2=NP1+NP1
NP3=NP1+NP2
DO 35 I=1,N2
II=2*I
I2=NP1-I
I1=NP3+II
A(I1-1)=TN1(I)
A(I1)=TN1(I2)
I1=NP2+II
A(I1-1)=0.0E0
A(I1)=0.0E0
I1=NP1+II
A(I1-1)=-TN2(I)
35 A(I1)=-TN2(I2+1)
A(NP2+2)=-TN2(1)
A(NP3-1)=-TN2(N2+1)
A(NP1+1)=0.0E0
A(NP1+2)=0.0E0
CALL BANDR(A(1),A(NP1+1),A(NP2+1),A(NP3+1),N)
CALL QL(A(1),A(NP1+1),N,IERR)
LINE=N*NEP1
J2=NP1
DO 40 I=1,N
LINE=LINE+1
J2=J2-1
40 R(LINE)=A(J2)
C IN THE FOLLOWING SEGMENT OF CODE, EIGENVALUES RELEVANT TO ALL
C GENERAL SEPARABLE PROBLEMS ARE CALCULATED.
50 DO 65 I=1,NEP1
I2=2*(I-1)
J2=N*(I-1)
DO 65 LNINDEX=1,NDP1,I2
LINE=ID(LNINDEX)
ISPAN=LNINDEX+I2
ISPAN=ID(ISPAN)
ISPAN=ISPAN-LINE-1
J1=N
I1=LINE
DO 55 J=1,ISPAN

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RTSC1750
RTSC1760
RTSC1770
RTSC1780
RTSC1790
RTSC1800
RTSC1810
RTSC1820
RTSC1830
RTSC1840
RTSC1850
RTSC1860
RTSC1870
RTSC1880
RTSC1890
RTSC1900
RTSC1910
RTSC1920
RTSC1930
RTSC1940
RTSC1950

RTSG0010
RTSG0020
RTSG0030
RTSG0040
RTSG0050
RTSG0060
RTSG0070
RTSG0080
RTSG0090
RTSG0100
RTSG0110
RTSG0120
RTSG0130
RTSG0140
RTSG0150
RTSG0160
RTSG0170
RTSG0180
RTSG0190
RTSG0200
RTSG0210
RTSG0220
RTSG0230
RTSG0240
RTSG0250
RTSG0260
RTSG0270
RTSG0280
RTSG0290
RTSG0300
RTSG0310
RTSG0320
RTSG0330
RTSG0340
RTSG0350
RTSG0360
RTSG0370
RTSG0380
RTSG0390
RTSG0400
RTSG0410
RTSG0420
RTSG0430
RTSG0440
RTSG0450
RTSG0460
RTSG0470
RTSG0480
RTSG0490
RTSG0500
RTSG0510
RTSG0520
RTSG0530

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      J1=J1+1
      I1=I1+1
      A(J)=TN1(I1)
55      A(J1)=-TN2(I1+1)
      A(J1)=0.0E0
      CALL QL(A(1),A(NP1),ISPAN,IERR)
      LINE=LINE+J2
      J1=ISPAN+1
      DO 60 J=1,ISPAN
          LINE=LINE+1
          J1=J1-1
60      R(LINE)=A(J1)
      R(LINE+1)=0.0E0
65      CONTINUE
      RETURN
      END
C
C
C
SUBROUTINE SORT(N,NTYPE,NE,ND,ID,R,A,IPI,KPI)
C      SUBROUTINE SORT ORDERS THE EIGENVALUES COMPUTED IN ROOTSC
C      OR ROOTSG ACCORDING TO SIZE. THIS MUST BE DONE TO INSURE THE
C      NUMERICAL STABILITY OF THE ALGORITHM. IN THE CASE OF CONSTANT
C      COEFFICIENT PROBLEMS, SORTING THE EIGENVALUES ALLOWS SOME
C      REDUCTION IN COMPUTATIONS TO TAKE PLACE.
C      THE ARRAY IPI KEEPS TRACK OF THE PERMUTATIONS.
C      VERSION DATE: FEBRUARY 1, 1977.
      REAL ID,IPI,KPI
      DIMENSION IPI(1),KPI(1),A(1),R(1),ID(1)
      IF(NE.EQ.0) RETURN
      NP1=N+1
      NN=NP1
      IF(NTYPE.GE.3) NN=N
      LINET=N*NE+1
      LINEB=LINET+NN-2
      RMIN=AMIN1(R(LINET),R(LINEB))-2.0E0
C      EPS IS ADDED TO ALL INTEGERS STORED IN THE REAL ARRAY IPI.
C      THIS PREVENTS UNFORTUNATE ROUNDING ERRORS FORM OCCURRING WHEN
C      THEY ARE CONVERTED BACK INTO INTEGERS.
      EPS=0.25E0
      DO 505 I=1,N
          A(I)=R(I)
          IPI(I)=FLOAT(I)+EPS
505      KPI(I)=IPI(I)
      DO 530 I=1,NE
          I2=2*I
          I1=I2/2
          DO 510 LNINDEX=I1,ND,I1
              LINE=ID(LNINDEX+1)
710      A(LINE)=RMIN
          A(NN)=RMIN
          DO 520 LNINDEX=I1,ND,I2
              LINET=LNINDEX-I1+1
              LINET=ID(LINET)
              LINET=LINET+1
              LINEB=ID(LNINDEX+1)
              LINEB=LINET+1
              LINE=LINET+1
              ISPAN=LNINDEX+I1+1
              ISPAN=ID(ISPAN)
              ISPAN=ISPAN-1
              KPI(LINET)=IPI(LINEB-1)
          DO 520 J=LINET,ISPAN
              IF(A(LINET).GT.A(LINEB)) GO TO 515
              KPI(J)=IPI(LINEB)
              LINEB=LINET+1
              GO TO 520
715      KPI(J)=IPI(LINET)
              LINET=LINET+1
720      CONTINUE
          LINE=N*I
          JPI=LINET
          DO 525 J=1,N
              IPI(J)=KPI(J)
              JPI=JPI+1
725      A(J)=R(JPI)

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RTSG0540
RTSG0550
RTSG0560
RTSG0570
RTSG0580
RTSG0590
RTSG0600
RTSG0610
RTSG0620
RTSG0630
RTSG0640
RTSG0650
RTSG0660
RTSG0670
RTSG0680
RTSG0690
RTSG0700

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SORT0010
SORT0020
SORT0030
SORT0040
SORT0050
SORT0060
SORT0070
SORT0080
SORT0090
SORT0100
SORT0110
SORT0120
SORT0130
SORT0140
SORT0150
SORT0160
SORT0170
SORT0180
SORT0190
SORT0200
SORT0210
SORT0220
SORT0230
SORT0240
SORT0250
SORT0260
SORT0270
SORT0280
SORT0290
SORT0300
SORT0310
SORT0320
SORT0330
SORT0340
SORT0350
SORT0360
SORT0370
SORT0380
SORT0390
SORT0400
SORT0410
SORT0420
SORT0430
SORT0440
SORT0450
SORT0460
SORT0470
SORT0480
SORT0490
SORT0500
SORT0510
SORT0520
SORT0530
SORT0540
SORT0550
SORT0560
SORT0570

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DO 530 J=1,N
  JPI=IPI(J)
  JPI=JPI+LINE
530 R(JPI)=A(J)
C THE FOLLOWING SEGMENT OF CODE SORTS THE SPECIAL EIGENVALUES
C WHICH ARISE IN THE CASE OF PERIODIC BOUNDARY CONDITIONS.
IF(NTYPE.LT.3) RETURN
LINE=N*(NE+1)
JPI=LINE
DO 540 J=1,N
  JPI=JPI+1
540 A(J)=R(JPI)
DO 545 J=1,N
  JPI=IPI(J)
  JPI=JPI+LINE
545 R(JPI)=A(J)
RETURN
END
C
C
SUBROUTINE QL(D,E,N,IERR)
C SUBROUTINE QL FINDS ALL THE EIGENVALUES OF A SYMMETRIC
C TRIDIAGONAL MATRIX, THE RATIONAL (SQUARE ROOT FREE) QL
C ALGORITHM IS USED. QL IS CALLED IN CONNECTION WITH FINDING
C EIGENVALUES FOR GENERAL SEPARABLE PROBLEMS, AND CONSTANT
C COEFFICIENT PROBLEMS WHERE ONE OR MORE MIXED BOUNDARY CONDITIONS
C ARE PRESENT. SUBROUTINE QL IS AN ADAPTATION OF SUBROUTINE
C TQLRAT FROM THE EISPACK SUBROUTINE LIBRARY.
C THE DIAGONAL OF THE TRIDIAGONAL MATRIX IS STORED IN THE ARRAY
C D, THE CODIAGONAL IN E. THE EIGENVALUES ARE WRITTEN IN D.
C VERSION DATE: MARCH 1, 1977.
COMMON /MACHEP/TOL
DIMENSION D(1),E(1)
IERR=0
IF(N.EQ.1) RETURN
NMI=N-1
DO 100 I=1,NMI
  E(I)=E(I)*E(I)
100 F=0.0E0
  B=0.0E0
  C=0.0E0
  E(N)=0.0E0
DO 290 L=1, N
  J=0
  H=TOL*(ABS(D(L))+SQRT(E(L)))
  IF(B.GT.H) GO TO 105
  B=H
  C=B*B
105 DO 110 M=L, N
  IF(E(M).LE.C) GO TO 120
110 CONTINUE
  M=N
120 IF(M.EQ.L) GO TO 210
130 IF(J.EQ.30) GO TO 1000
  J=J+1
  L1=L+1
  S=SQRT(E(L))
  G=D(L)
  P=(D(L1)-G)/(2.0E0*S)
  R=SQRT(P*P+1.0E0)
  D(L)=S/(P+SIGN(R,P))
  H=G-D(L)
DO 140 I=L1, N
  D(I)=D(I)-H
140 F=F+H
  G=D(M)
  IF(G.EQ.0.0E0) G=B
  H=G
  S=0.0E0
  MML=M-L
DO 200 II=1, MML
  I=M-II
  P=C*H
  R=P+E(I)
  E(I+1)=S*R
  S=E(I)/R

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SORT0580
SORT0590
SORT0600
SORT0610
SORT0620
SORT0630
SORT0640
SORT0650
SORT0660
SORT0670
SORT0680
SORT0690
SORT0700
SORT0710
SORT0720
SORT0730
SORT0740
SORT0750
QL000010
QL000020
QL000030
QL000040
QL000050
QL000060
QL000070
QL000080
QL000090
QL000100
QL000110
QL000120
QL000130
QL000140
QL000150
QL000160
QL000170
QL000180
QL000190
QL000200
QL000210
QL000220
QL000230
QL000240
QL000250
QL000260
QL000270
QL000280
QL000290
QL000300
QL000310
QL000320
QL000330
QL000340
QL000350
QL000360
QL000370
QL000380
QL000390
QL000400
QL000410
QL000420
QL000430
QL000440
QL000450
QL000460
QL000470
QL000480
QL000490
QL000500
QL000510
QL000520
QL000530
QL000540
QL000550
QL000560

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                D(I+1)=H+S*(H+D(I))
                G=D(I)-E(I)/G
                IF(G.EQ.0.0E0) G=B
                H=G*P/R
200             CONTINUE
                E(L)=S*G
                D(L)=H
                IF(H.EQ.0.0E0) GO TO 210
                IF(ABS(E(L)).LE.ABS(C/H)) GO TO 210
                E(L)=H*E(L)
                IF(E(L).NE.0.0E0) GO TO 130
210             P=D(L)+F
                IF(L.EQ.1) GO TO 250
                DO 230 II=2, L
                    I=L+2-II
                    IF(P.GE.D(I-1)) GO TO 270
                    D(I)=D(I-1)
230             CONTINUE
250             I=1
270             D(I)=P
290             CONTINUE
                GO TO 1001
1000            IERR=L
1001            RETURN
                END
C
C
C             SUBROUTINE BANDR(D,E,W,V,N)
C             SUBROUTINE BANDR IS USED TO REDUCE A SYMMETRIC PENTA-
C             DIAGONAL MATRIX TO A SYMMETRIC TRIDIAGONAL MATRIX. BANDR IS
C             ENTERED ONLY IN THE CASE NTYPE = 3, AND IS CALLED IN
C             CONNECTION WITH FINDING THE SPECIAL EIGENVALUES ASSOCIATED
C             WITH THE PERIODIC BOUNDARY CONDITIONS. BANDR IS AN ADAPTATION
C             OF SUBROUTINE BANDR FROM THE EISPACK SUBROUTINE LIBRARY.
C             THE DIAGONAL OF THE PENTADIAGONAL MATRIX IS STORED IN V,
C             THE FIRST CO-DIAGONAL IN W, AND THE SECOND CO-DIAGONAL IN E.
C             THE DIAGONAL OF THE TRIDIAGONAL MATRIX IS WRITTEN IN D, AND
C             THE CO-DIAGONAL IS WRITTEN IN E.
C             VERSION DATE: FEBRUARY 1, 1977.
C             DIMENSION D(1),E(1),V(1),W(1)
                DO 30 J=1, N
30             D(J)=1.0E0
                NM2=N-2
                DO 700 K=1, NM2
                    KP2=K+2
                    G=E(KP2)
                    E(K+1)=W(K+1)
                    DO 500 J=KP2, N, 2
                        JM1=J-1
                        JP1=J+1
                        JP2=J+2
                        IF (G.EQ.0.0E0) GO TO 700
                        B1=E(JM1)/G
                        B2=B1*D(JM1)/D(J)
                        S2=1.0E0/(1.0E0+B1*B2)
                        IF (S2.GE.0.5E0) GO TO 450
                        B1=G/E(JM1)
                        B2=B1*D(J)/D(JM1)
                        C2=1.0E0-S2
                        D(JM1)=C2*D(JM1)
                        D(J)=C2*D(J)
                        F1=2.0E0*W(J)
                        F2=B1*V(JM1)
                        W(J)=-B2*(B1*W(J)-V(J))-F2+W(J)
                        V(JM1)=B2*(B2*V(J)+F1)+V(JM1)
                        V(J)=B1*(F2-F1)+V(J)
                        U=W(JM1)+B2*E(J)
                        E(J)=-B1*W(JM1)+E(J)
                        W(JM1)=U
                        E(JM1)=E(JM1)+B2*G
                        IF (J.EQ.N) GO TO 500
                        U=E(JP1)+B2*W(JP1)
                        W(JP1)=-B1*E(JP1)+W(JP1)
                        E(JP1)=U
                        IF (JP2.GT.N) GO TO 500
                        G=B2*E(JP2)
BNDR0010
BNDR0020
BNDR0030
BNDR0040
BNDR0050
BNDR0060
BNDR0070
BNDR0080
BNDR0090
BNDR0100
BNDR0110
BNDR0120
BNDR0130
BNDR0140
BNDR0150
BNDR0160
BNDR0170
BNDR0180
BNDR0190
BNDR0200
BNDR0210
BNDR0220
BNDR0230
BNDR0240
BNDR0250
BNDR0260
BNDR0270
BNDR0280
BNDR0290
BNDR0300
BNDR0310
BNDR0320
BNDR0330
BNDR0340
BNDR0350
BNDR0360
BNDR0370
BNDR0380
BNDR0390
BNDR0400
BNDR0410
BNDR0420
BNDR0430
BNDR0440
BNDR0450
BNDR0460
BNDR0470
BNDR0480
BNDR0490

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450      GO TO 500
        U=D(JM1)
        D(JM1)=S2*D(J)
        D(J)=S2*U
        F1=2.0E0*W(J)
        F2=B1*V(J)
        U=B1*(F2-F1)+V(JM1)
        W(J)=B2*(B1*W(J)-V(JM1))+F2-W(J)
        V(JM1)=B2*(B2*V(JM1)+F1)+V(J)
        V(J)=U
        U=B2*W(JM1)+E(J)
        E(J)=-W(JM1)+B1*E(J)
        W(JM1)=U
        E(JM1)=B2*E(JM1)+G
        IF (J.EQ.N) GO TO 500
        U=B2*E(JP1)+W(JP1)
        W(JP1)=-E(JP1)+B1*W(JP1)
        E(JP1)=U
        IF (JP2.GT.N) GO TO 500
        G=E(JP2)
        E(JP2)=B1*E(JP2)
500      CONTINUE
700      U=1.0E0
        DO 850 J=2, N
            E(J)=SQRT(D(J))
            D(J)=D(J)*V(J)
            E(J-1)=U*E(J)*W(J)
850      U=E(J)
        D(1)=V(1)
        E(N)=0.0E0
        RETURN
        END
C
C
        SUBROUTINE STEP1(N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2,
1          IBDIM,B,ND,ID,R,A)
C          SUBROUTINE STEP1 CARRIES OUT THE INITIAL MARCHING PHASE OF
C          THE GENERALIZED MARCHING ALGORITHM. THE SYSTEM OF EQUATIONS
C          IS REDUCED TO AN EQUIVALENT SET FOR THE UNKNOWNNS ON THE ND
C          SEPARATING LINES.
C          VERSION DATE: MARCH 1, 1977.
        REAL ID
        DIMENSION TN1(1),TN2(1),TM1(1),TM2(1),A(1),B(1),R(1),ID(1)
        NN=ND
        IF(NTYPE.GE.3) NN=ND+1
        IF(NN.LT.1) RETURN
        M1=M+1
        M2=M1+M
        ITYPE=NTYPE-(NTYPE/2)*2
        ID2=ID(2)
        DO 130 LNINDX=1,NN
            LINE=ID(LNINDX+1)
            LINEA=ID(LNINDX+2)
            LINEB=ID(LNINDX)
            LINEC=LINE
            IF(LINE.EQ.N) LINEC=0
            ISPANA=LINEA-LINE-1
            IF(ISPANA.LE.0) ISPANA=ID2-1
            ISPANB=LINE-LINEB-1
C          CHECK FOR POSSIBLE REDUCTION IN COMPUTATION IN THE
C          CONSTANT COEFFICIENT CASE.
            IF((ITYPE.EQ.1).OR.(ISPANA.NE.ISPANB)) GO TO 105
            IF((TN1(LINEB+1).EQ.TN1(LINEA-1)).AND.
1          (TN2(LINEB+2).EQ.TN2(LINEA-1))) GO TO 120
C          THE FOLLOWING CODE IS EXECUTED WHEN NO REDUCTION
C          IN COMPUTATION IS POSSIBLE.
105      CALL MARCH1(M,MTYPE,TM1,TM2,TN1,TN2,IBDIM,B,
1          ISPANB,LINE,-1,A(1),A(M1))
            INDXRT=LINEB+1
            CALL TRI1(M,MTYPE,TM1,TM2,ISPANB,R(INDXRT),TN2(INDXRT),
1          A(1),A(M1),A(M2))
            T1=TN2(LINE)
            LNPTR=(LINE-1)*IBDIM
            DO 110 ICOMP=1,M
                LNPTR=LNPTR+1

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BNDR0500
BNDR0510
BNDR0520
BNDR0530
BNDR0540
BNDR0550
BNDR0560
BNDR0570
BNDR0580
BNDR0590
BNDR0600
BNDR0610
BNDR0620
BNDR0630
BNDR0640
BNDR0650
BNDR0660
BNDR0670
BNDR0680
BNDR0690
BNDR0700
BNDR0710
BNDR0720
BNDR0730
BNDR0740
BNDR0750
BNDR0760
BNDR0770
BNDR0780
BNDR0790
BNDR0800
BNDR0810
BNDR0820

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STP10010
STP10020
STP10030
STP10040
STP10050
STP10060
STP10070
STP10080
STP10090
STP10100
STP10110
STP10120
STP10130
STP10140
STP10150
STP10160
STP10170
STP10180
STP10190
STP10200
STP10210
STP10220
STP10230
STP10240
STP10250
STP10260
STP10270
STP10280
STP10290
STP10300
STP10310
STP10320
STP10330
STP10340
STP10350
STP10360
STP10370
STP10380
STP10390
STP10400
STP10410

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110      B(LNPTR)=B(LNPTR)-T1*A(ICOMP)          STP10420
      CALL MARCH1(M,MTYPE,TM1,TM2,TN1,TN2,IBDIM,B, STP10430
1      ISPAN,LINEC,1,A(1),A(M1))             STP10440
      INDXRT=LINEC+1                          STP10450
      CALL TRI1(M,MTYPE,TM1,TM2,ISPANA,R(INDXRT),TN2(INDXRT), STP10460
1      A(1),A(M1),A(M2))                    STP10470
      INDXRT=INDXRT+ISPANA                    STP10480
      T1=TN2(INDXRT)                          STP10490
      LNPTR=(LINE-1)*IBDIM                    STP10500
      DO 115 ICOMP=1,M                        STP10510
          LNPTR=LNPTR+1                       STP10520
115      B(LNPTR)=B(LNPTR)-T1*A(ICOMP)        STP10530
      GO TO 130                               STP10540
C      THE FOLLOWING CODE IS EXECUTED FOR CONSTANT COEFFICIENT STP10550
C      PROBLEMS IN WHICH A REDUCTION IN COMPUTATION IS POSSIBLE. STP10560
120     CALL MARCH3(M,MTYPE,TM1,TM2,TN1,TN2,IBDIM,B, STP10570
1     ISPANB,LINE.LINEC,A(1),A(M1))          STP10580
      INDXRT=LINEB+1                          STP10590
      CALL TRI1(M,MTYPE,TM1,TM2,ISPANB,R(INDXRT),TN2(INDXRT), STP10600
1     A(1),A(M1),A(M2))                    STP10610
      T1=TN2(LINE)                            STP10620
      LNPTR=(LINE-1)*IBDIM                    STP10630
      DO 125 ICOMP=1,M                        STP10640
          LNPTR=LNPTR+1                       STP10650
125      B(LNPTR)=B(LNPTR)-T1*A(ICOMP)        STP10660
130     CONTINUE                             STP10670
      RETURN                                  STP10680
      END                                     STP10690

C
C
C      SUBROUTINE STEP2(N,NTYPE,TN2,M,MTYPE,TM1,TM2, STP20010
1      IBDIM,B,NE,ND,ID,R,A)                 STP20020
C      SUBROUTINE STEP2 USES NE-1 2-REDUCTION STEPS TO REDUCE STP20030
C      THE SET OF EQUATIONS FOR THE ND SEPARATING LINES TO A MATRIX STP20040
C      EQUATION FOR A SINGLE LINE OF UNKNOWNNS. APPROXIMATELY STP20050
C      HALF OF THE REMAINING LINES ARE ELIMINATED AT EACH RECURSIVE STP20060
C      STEP.                                  STP20070
C      VERSION DATE: MARCH 1, 1977.           STP20080
      REAL ID                                  STP20090
      DIMENSION TN2(1),TM1(1),TM2(1),A(1),B(1),R(1),ID(1) STP20100
      NN=NE-1                                  STP20110
      IF(NTYPE.GE.3) NN=NE                     STP20120
      IF(NN.LT.1) RETURN                       STP20130
      I=M+1                                     STP20140
      M2=M1+M                                  STP20150
      M3=M2+M                                  STP20160
      ITYPE=NTYPE-(NTYPE/2)*2                  STP20170
      DO 240 I=1,NN                             STP20180
          I2=2*I                                 STP20190
          I1=I2/2                               STP20200
          J1=(I-1)*N+1                          STP20210
          J2=I*N+1                              STP20220
      DO 240 LNINDX=I1,ND,I2                    STP20230
          LINE=ID(LNINDX+1)                     STP20240
          LINEA=LNINDX+1+I1                      STP20250
          LINEA=ID(LINEA)                       STP20260
          LINEB=LNINDX+1-I1                      STP20270
          LINEB=ID(LINEB)                       STP20280
          ISPANB=LINEA-LINE                     STP20290
          ISPANB=LINE-LINEB                     STP20300
          IF((LINEB.EQ.0).AND.(NTYPE.LT.3)) GO TO 215 STP20310
          LNPTR=(LINE-1)*IBDIM                   STP20320
          DO 205 ICOMP=1,M                       STP20330
              LNPTR=LNPTR+1                     STP20340
205      A(ICOMP)=B(LNPTR)                     STP20350
          INDXRT=LINEB+J2                       STP20360
          INDXTN=LINEB+1                       STP20370
          CALL TRI1(M,MTYPE,TM1,TM2,ISPANB,R(INDXRT),TN2(INDXTN), STP20380
1          A(1),A(M1),A(M2))                   STP20390
          IRTDIF=LINE+J1                       STP20400
          INDXRT=INDXRT+ISPANB                  STP20410
          ISPAN=ISPANB-1                        STP20420
          CALL TRI2(M,MTYPE,TM1,TM2,ISPAN,R(INDXRT),R(IRTDIF), STP20430
1          A(1),A(M1),A(M2),A(M3))           STP20440
          ISPAN=LINEB                           STP20450
          IF(LINEB.EQ.0) ISPAN=N                STP20460

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                LNPTR=(ISPAN-1)*IBDIM
                DO 210 ICOMP=1,M
                LNPTR=LNPTR+1
210             B(LNPTR)=B(LNPTR)+A(ICOMP)
C             CHECK FOR POSSIBLE REDUCTION IN COMPUTATION IN THE
C             CONSTANT COEFFICIENT CASE.
                IF(LINEA.GT.N) GO TO 240
                IF((ITYPE.EQ.0).AND.(ISPANA.EQ.ISPANB)) GO TO 225
C             THE FOLLOWING CODE IS EXECUTED WHEN NO REDUCTION
C             IN COMPUTATION IS POSSIBLE.
215             LNPTR=(LINE-1)*IBDIM
                DO 220 ICOMP=1,M
                LNPTR=LNPTR+1
220             A(ICOMP)=B(LNPTR)
                INDXRT=LINEB+J2
                INDXTN=LINE+1
                CALL TRI1(M,MTYPE,TM1,TM2,ISPANA,R(INDXRT),TN2(INDXTN),
                1             A(1),A(M1),A(M2))
                INDXRT=INDXRT+ISPANA
                IRTDIF=LINEB+J1
                ISPAN=ISPANB-1
                CALL TRI2(M,MTYPE,TM1,TM2,ISPAN,R(INDXRT),R(IRTDIF),
                1             A(1),A(M1),A(M2),A(M3))
225             LNPTR=(LINEA-1)*IBDIM
                DO 230 ICOMP=1,M
                LNPTR=LNPTR+1
230             B(LNPTR)=B(LNPTR)+A(ICOMP)
240             CONTINUE
                RETURN
                END

                SUBROUTINE STEP3(N,NTYPE,TN2,M,MTYPE,TM1,TM2,
                1             IBDIM,B,NE,ND,ID,R,A)
C             IN SUBROUTINE STEP3, THE UNKNOWNNS ON THE ND SEPARATING
C             LINES ARE DETERMINED. THE LINES ARE DETERMINED IN THE
C             REVERSE ORDER THAT THEY WERE ELIMINATED IN STEP2.
C             VERSION DATE: MARCH 1, 1977.
                REAL ID
                DIMENSION TN2(1),TM1(1),TM2(1),A(1),B(1),R(1),ID(1)
                NP1=N+1
                M1=M+1
                M2=M1+M
                M3=M2+M
                ITYPE=NTYPE-(NTYPE/2)*2
C             IN THE FOLLOWING SEGMENT OF CODE, THE SPECIAL LINE OF UNKNOWNNS
C             ASSOCIATED WITH PERIODIC BOUNDARY CONDITIONS IS DETERMINED.
250             IF(NTYPE.LT.3) GO TO 270
                NM1=N-1
                LNPTR=NM1*IBDIM
                DO 255 ICOMP=1,M
                LNPTR=LNPTR+1
255             A(ICOMP)=B(LNPTR)
                INDXRT=(NE+1)*N+1
                IRTDIF=NE*N+1
                CALL TRI2(M,MTYPE,TM1,TM2,NM1,R(INDXRT),R(IRTDIF),
                1             A(1),A(M1),A(M2),A(M3))
                INDXRT=INDXRT+NM1
                CALL TRI1(M,MTYPE,TM1,TM2,1,R(INDXRT),TN2(NP1),
                1             A(1),A(M1),A(M2))
                LNPTR=NM1*IBDIM
                L2=LNPTR-IBDIM
                T1=TN2(1)
                T2=TN2(N)
                DO 265 ICOMP=1,M
                LNPTR=LNPTR+1
                L2=L2+1
                X=A(ICOMP)
                B(LNPTR)=X
                B(L2)=B(L2)+T2*X
265             B(ICOMP)=B(ICOMP)+T1*X
270             IF(NE.EQ.0) RETURN
C             IN THE FOLLOWING SEGMENT OF CODE, THE UNKNOWNNS ON THE
C             REMAINING SEPARATING LINES ARE DETERMINED.
                DO 345 II=1,NE
                I=NE+1-II
                STP20470
                STP20480
                STP20490
                STP20500
                STP20510
                STP20520
                STP20530
                STP20540
                STP20550
                STP20560
                STP20570
                STP20580
                STP20590
                STP20600
                STP20610
                STP20620
                STP20630
                STP20640
                STP20650
                STP20660
                STP20670
                STP20680
                STP20690
                STP20700
                STP20710
                STP20720
                STP20730
                STP20740
                STP20750
                STP20760

                STP30010
                STP30020
                STP30030
                STP30040
                STP30050
                STP30060
                STP30070
                STP30080
                STP30090
                STP30100
                STP30110
                STP30120
                STP30130
                STP30140
                STP30150
                STP30160
                STP30170
                STP30180
                STP30190
                STP30200
                STP30210
                STP30220
                STP30230
                STP30240
                STP30250
                STP30260
                STP30270
                STP30280
                STP30290
                STP30300
                STP30310
                STP30320
                STP30330
                STP30340
                STP30350
                STP30360
                STP30370
                STP30380
                STP30390
                STP30400
                STP30410
                STP30420
                STP30430
                STP30440

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      I2=2**I
      I1=I2/2
      J1=(I-1)*N+1
      J2=I*N+1
DO 345 LNINDX=I1,ND,I2
      LINE=ID(LNINDX+1)
      LINEA=LNINDX+1+I1
      LINEA=ID(LINEA)
      LINEB=LNINDX+1-I1
      LINEB=ID(LINEB)
      ISPANA=LINEA-LINE-1
      ISPANB=LINE-LINEB-1
C      IN THE FOLLOWING SEGMENT OF CODE, THE RIGHT HAND SIDE IS
C      MODIFIED USING UNKNOWNNS WHICH WERE SOLVED FOR ON A PREVIOUS
C      STEP OF THE RECURSION.
      ISPAN=LINEB
      IF(LINEB.EQ.0) ISPAN=ID(ND+2)
      IF(ISPAN.GT.N) GO TO 295
      LNPTR=(ISPAN-1)*IBDIM
      DO 275 ICOMP=1,M
          LNPTR=LNPTR+1
275      A(ICOMP)=B(LNPTR)
C      CHECK FOR POSSIBLE REDUCTION IN COMPUTATION IN THE
C      CONSTANT COEFFICIENT CASE.
      IF((ITYPE.EQ.1).OR.(LINEA.GT.N)) GO TO 285
      IF(ISPANA.NE.ISPANB) GO TO 285
      LNPTR=(LINEA-1)*IBDIM
      DO 280 ICOMP=1,M
          LNPTR=LNPTR+1
280      A(ICOMP)=A(ICOMP)+B(LNPTR)
      GO TO 305
C      THE FOLLOWING CODE IS EXECUTED WHEN NO REDUCTION
C      IN COMPUTATION IS POSSIBLE.
285      INDXRT=LINEB+J1
      INDXTN=LINEB+1
      CALL TRI1(M,MTYPE,TM1,TM2,ISPANB,R(INDXRT),TN2(INDXTN),
1          A(1),A(M1),A(M2))
      T1=TN2(LINE)
      LNPTR=(LINE-1)*IBDIM
      DO 290 ICOMP=1,M
          LNPTR=LNPTR+1
290      B(LNPTR)=B(LNPTR)+T1*A(ICOMP)
295      IF(LINEA.GT.N) GO TO 315
      LNPTR=(LINEA-1)*IBDIM
      DO 300 ICOMP=1,M
          LNPTR=LNPTR+1
300      A(ICOMP)=B(LNPTR)
305      INDXRT=LINE+J1
      INDXTN=LINE+1
      CALL TRI1(M,MTYPE,TM1,TM2,ISPANA,R(INDXRT),TN2(INDXTN),
1          A(1),A(M1),A(M2))
      T1=TN2(LINEA)
      LNPTR=(LINE-1)*IBDIM
      DO 310 ICOMP=1,M
          LNPTR=LNPTR+1
310      B(LNPTR)=B(LNPTR)+T1*A(ICOMP)
C      IN THE FOLLOWING SEGMENT OF CODE, WE SOLVE FOR THE NEW LINE
C      OF UNKNOWNNS, AND THEN MODIFY THE RIGHT HAND SIDES OF THE
C      ND+1 SMALLER BLOCKS IN THE APPROPRIATE FASHION.
315      LNPTR=(LINE-1)*IBDIM
      DO 330 ICOMP=1,M
          LNPTR=LNPTR+1
330      A(ICOMP)=B(LNPTR)
      INDXRT=LINEB+J2
      IRTDIF=LINEB+J1
      CALL TRI2(M,MTYPE,TM1,TM2,ISPANB,R(INDXRT),R(IRTDIF),
1          A(1),A(M1),A(M2),A(M3))
      INDXRT=LINE+J2
      IRTDIF=LINE+J1
      CALL TRI1(M,MTYPE,TM1,TM2,1,R(INDXRT-1),TN2(NP1),
1          A(1),A(M1),A(M2))
      CALL TRI2(M,MTYPE,TM1,TM2,ISPANA,R(INDXRT),R(IRTDIF),
1          A(1),A(M1),A(M2),A(M3))
      LNPTR=(LINE-1)*IBDIM
      L2=LNPTR-IBDIM
      L1=LNPTR+IBDIM
      STP30450
      STP30460
      STP30470
      STP30480
      STP30490
      STP30500
      STP30510
      STP30520
      STP30530
      STP30540
      STP30550
      STP30560
      STP30570
      STP30580
      STP30590
      STP30600
      STP30610
      STP30620
      STP30630
      STP30640
      STP30650
      STP30660
      STP30670
      STP30680
      STP30690
      STP30700
      STP30710
      STP30720
      STP30730
      STP30740
      STP30750
      STP30760
      STP30770
      STP30780
      STP30790
      STP30800
      STP30810
      STP30820
      STP30830
      STP30840
      STP30850
      STP30860
      STP30870
      STP30880
      STP30890
      STP30900
      STP30910
      STP30920
      STP30930
      STP30940
      STP30950
      STP30960
      STP30970
      STP30980
      STP30990
      STP31000
      STP31010
      STP31020
      STP31030
      STP31040
      STP31050
      STP31060
      STP31070
      STP31080
      STP31090
      STP31100
      STP31110
      STP31120
      STP31130
      STP31140
      STP31150
      STP31160
      STP31170
      STP31180
      STP31190
      STP31200

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      T1=TN2(LINE+1)
      T2=TN2(LINE)
      DO 345 ICOMP=1,M
        LNPTR=LNPTR+1
        L1=L1+1
        L2=L2+1
        X=A(ICOMP)
        B(LNPTR)=X
        B(L1)=B(L1)+T1*X
        B(L2)=B(L2)+T2*X
345      RETURN
      END
C
C
      SUBROUTINE STEP4(TN1,TN2,M,MTYPE,TM1,TM2,IBDIM,B,ND,ID,R,A)
C      IN SUBROUTINE STEP4, THE ND+1 SMALLER PROBLEMS ARE
C      SOLVED USING THE MARCHING ALGORITHM. THIS COMPLETES
C      THE BACKSOLUTION PHASE OF THE GENERALIZED MARCHING ALGORITHM.
C      VERSION DATE: MARCH 1, 1977.
      REAL ID
      DIMENSION TN1(1),TN2(1),TM1(1),TM2(1),A(1),B(1),R(1),ID(1)
      M1=M+1
      M2=M1+M
      NDP1=ND+1
      DO 440 LNINDEX=1,NDP1
        LINE=ID(LNINDEX)
        ISPAN=ID(LNINDEX+1)
        ISPAN=ISPAN-LINE-1
        CALL MARCH1(M,MTYPE,TM1,TM2,TN1,TN2,IBDIM,B,
1         ISPAN,LINE,1,A(1),A(M1))
        INDXRT=LINE+1
        INDXTN=INDXRT+1
        CALL TRI1(M,MTYPE,TM1,TM2,ISPAN,R(INDXRT),TN2(INDXTN),
1         A(1),A(M1),A(M2))
        CALL MARCH2(M,MTYPE,TM1,TM2,TN1,TN2,IBDIM,B,
1         ISPAN,LINE,A(1),A(M1))
440      CONTINUE
      RETURN
      END
C
C
      SUBROUTINE TRI1(M,MTYPE,TM1,TM2,ISPAN,R,S,V,U,E)
C      SUBROUTINE TRI1 SOLVES A LINEAR SYSTEM WHICH INVOLVES
C      A POLYNOMIAL OF DEGREE ISPAN IN THE MATRIX TM. THE ZEROES OF
C      THE POLYNOMIAL ARE A SUBSET OF THE SET OF EIGENVALUES
C      COMPUTED DURING THE PREPROCESSING PHASE. THE LEADING COEFFICIENT
C      IS GIVEN AS A PRODUCT OF ISPAN SCALARS STORED IN S.
C      THE RIGHT HAND SIDE IS STORED IN V. THE SOLUTION IS COMPUTED
C      BY SOLVING ISPAN LINEAR SYSTEMS INVOLVING THE FACTORS OF THE
C      POLYNOMIAL. FOUR COMPLETE LOOPS ARE INCLUDED; THE CORRECT LOOP
C      FOR A GIVEN PROBLEM IS DETERMINED BY THE VALUE OF MTYPE.
C      VERSION DATE: FEBRUARY 1, 1977.
      DIMENSION TM1(1),TM2(1),U(1),V(1),E(1),R(1),S(1)
      MM1=M-1
      GO TO (3,50,20,75),MTYPE
C      THE LOOP FOR MTYPE = 1.
3      DO 15 J=1,ISPAN
        D=R(J)
        DS=S(J)
        U(1)=TM1(1)+D
        DO 5 ICOMP=2,M
          ICM1=ICOMP-1
          Q=TM2(ICOMP)/U(ICM1)
          V(ICOMP)=V(ICOMP)+V(ICM1)*Q
5          U(ICOMP)=TM1(ICOMP)+D-Q*TM2(ICOMP)
          Q=V(M)/U(M)
          V(M)=Q*DS
          DO 10 II=1,MM1
            ICOMP=M-II
            Q=(V(ICOMP)+TM2(ICOMP+1)*Q)/U(ICOMP)
10          V(ICOMP)=Q*DS
15      CONTINUE
      RETURN
C      THE LOOP FOR MTYPE = 3.
20     DO 40 J=1,ISPAN
        D=R(J)

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STP31210
STP31220
STP31230
STP31240
STP31250
STP31260
STP31270
STP31280
STP31290
STP31300
STP31310
STP31320

STP40010
STP40020
STP40030
STP40040
STP40050
STP40060
STP40070
STP40080
STP40090
STP40100
STP40110
STP40120
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STP40140
STP40150
STP40160
STP40170
STP40180
STP40190
STP40200
STP40210
STP40220
STP40230
STP40240
STP40250

TRI10010
TRI10020
TRI10030
TRI10040
TRI10050
TRI10060
TRI10070
TRI10080
TRI10090
TRI10100
TRI10110
TRI10120
TRI10130
TRI10140
TRI10150
TRI10160
TRI10170
TRI10180
TRI10190
TRI10200
TRI10210
TRI10220
TRI10230
TRI10240
TRI10250
TRI10260
TRI10270
TRI10280
TRI10290
TRI10300
TRI10310
TRI10320
TRI10330
TRI10340
TRI10350

	DS=S(J)	TRI10360
	U(1)=TM1(1)+D	TRI10370
	E(1)=TM2(1)	TRI10380
	DO 25 ICOMP=2,MM1	TRI10390
	ICM1=ICOMP-1	TRI10400
	Q=TM2(ICOMP)/U(ICM1)	TRI10410
	V(ICOMP)=V(ICOMP)+V(ICM1)*Q	TRI10420
	U(ICOMP)=TM1(ICOMP)+D-Q*TM2(ICOMP)	TRI10430
25	E(ICOMP)=Q*E(ICM1)	TRI10440
	E(MM1)=(E(MM1)+TM2(M))/U(MM1)	TRI10450
	V(MM1)=V(MM1)/U(MM1)	TRI10460
	DO 30 II=2,MM1	TRI10470
	ICOMP=M-II	TRI10480
	ICP1=ICOMP+1	TRI10490
	V(ICOMP)=(V(ICOMP)+TM2(ICP1)*V(ICP1))/U(ICOMP)	TRI10500
30	E(ICOMP)=(E(ICOMP)+TM2(ICP1)*E(ICP1))/U(ICOMP)	TRI10510
	Q=TM1(M)+D-TM2(1)*E(1)-TM2(M)*E(MM1)	TRI10520
	Q=(V(M)+TM2(1)*V(1)+TM2(M)*V(MM1))/Q	TRI10530
	V(M)=Q*DS	TRI10540
	DO 35 ICOMP=1,MM1	TRI10550
35	V(ICOMP)=(V(ICOMP)+E(ICOMP)*Q)*DS	TRI10560
40	CONTINUE	TRI10570
	RETURN	TRI10580
C	THE LOOP FOR MTYPE = 2.	TRI10590
50	MM2=M-2	TRI10600
	T3=1.0E0/TM2(3)	TRI10610
	T2=TM2(2)*T3	TRI10620
	T4=TM2(M)*T3	TRI10630
	DO 65 J=1,ISPAN	TRI10640
	D=R(J)	TRI10650
	DS=S(J)*T3	TRI10660
	U(1)=TM2(3)/(TM1(1)+D)	TRI10670
	T1=(TM1(2)+D)*T3	TRI10680
	Q=T2*U(1)	TRI10690
	V(2)=V(2)+V(1)*Q	TRI10700
	U(2)=1.0E0/(T1-Q*T2)	TRI10710
	DO 55 ICOMP=3,MM1	TRI10720
	ICM1=ICOMP-1	TRI10730
	V(ICOMP)=V(ICOMP)+V(ICM1)*U(ICM1)	TRI10740
55	U(ICOMP)=1.0E0/(T1-U(ICM1))	TRI10750
	Q=T4*U(MM1)	TRI10760
	V(M)=V(M)+V(MM1)*Q	TRI10770
	U(M)=(TM1(M)+D)*T3-Q*T4	TRI10780
	Q=V(M)/U(M)	TRI10790
	V(M)=Q*DS	TRI10800
	Q=(V(MM1)+Q*T4)*U(MM1)	TRI10810
	V(MM1)=Q*DS	TRI10820
	DO 60 II=2,MM2	TRI10830
	ICOMP=M-II	TRI10840
	Q=(V(ICOMP)+Q)*U(ICOMP)	TRI10850
60	V(ICOMP)=Q*DS	TRI10860
	Q=(V(1)+T2*Q)*U(1)	TRI10870
	V(1)=Q*DS	TRI10880
65	CONTINUE	TRI10890
	RETURN	TRI10900
C	THE LOOP FOR MTYPE = 4.	TRI10910
75	T2=1.0E0/TM2(2)	TRI10920
	M1=MM1/2	TRI10930
	M1P1=M1+1	TRI10940
	M1P2=M1+2	TRI10950
	M1P3=M1+3	TRI10960
	MSW=M-(M/2)*2	TRI10970
	M2=M1-MSW	TRI10980
	RM=1.0E0	TRI10990
	RA=1.0E0	TRI11000
	IF(MSW.EQ.1) GO TO 80	TRI11010
	RM=2.0E0	TRI11020
	RA=0.0E0	TRI11030
80	DO 85 ICOMP=1,M1	TRI11040
	JCOMP=M+MSW-ICOMP	TRI11050
	Q=(V(JCOMP)-V(ICOMP))/2.0E0	TRI11060
	V(JCOMP)=(V(JCOMP)+V(ICOMP))/2.0E0	TRI11070
85	V(ICOMP)=Q	TRI11080
	DO 115 J=1,ISPAN	TRI11090
	D=R(J)	TRI11100
	DS=S(J)*T2	TRI11110

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      T1=(TM1(1)+D)*T2
      U(1)=1.0E0/(T1+RA)
      DO 90 ICOMP=2,M1
          ICM1=ICOMP-1
          V(ICOMP)=V(ICOMP)+V(ICM1)*U(ICM1)
      90 U(ICOMP)=1.0E0/(T1-U(ICM1))
      U(M1P1)=1.0E0/T1
      V(M1P2)=V(M1P2)+V(M1P1)*U(M1P1)
      U(M1P2)=1.0E0/(T1-2.0E0*U(M1P1))
      DO 95 ICOMP=M1P3,MM1
          ICM1=ICOMP-1
          V(ICOMP)=V(ICOMP)+V(ICM1)*U(ICM1)
      95 U(ICOMP)=1.0E0/(T1-U(ICM1))
      V(M)=V(M)+V(MM1)*U(MM1)*RM
      U(M)=T1-RM*U(MM1)-RA
      Q=V(M)/U(M)
      V(M)=Q*DS
      DO 105 II=1,M2
          ICOMP=M-II
          Q=(V(ICOMP)+Q)*U(ICOMP)
      105 V(ICOMP)=Q*DS
      Q=(V(M1P1)+2.0E0*Q)*U(M1P1)
      V(M1P1)=Q*DS
      Q=V(M1)*U(M1)
      V(M1)=Q*DS
      DO 110 II=2,M1
          ICOMP=M1P1-II
          Q=(V(ICOMP)+Q)*U(ICOMP)
      110 V(ICOMP)=Q*DS
      115 CONTINUE
      DO 120 ICOMP=1,M1
          JCOMP=M+MSW-ICOMP
          Q=V(JCOMP)-V(ICOMP)
          V(JCOMP)=V(JCOMP)+V(ICOMP)
      120 V(ICOMP)=Q
      RETURN
      END

      SUBROUTINE TRI2(M,MTYPE,TM1,TM2,ISPAN,R1,R2,V,W,U,E)
      C SUBROUTINE TRI2 SOLVES A LINEAR SYSTEM WHICH INVOLVES
      C A RATIONAL FUNCTION OF THE MATRIX TM. THE DEGREE OF BOTH THE
      C NUMERATOR AND THE DENOMINATOR IS ISPAN. THE ZEROES OF BOTH
      C POLYNOMIALS ARE SUBSETS OF THE EIGENVALUES CALCULATED DURING
      C THE PREPROCESSING PHASE. THE ZEROES OF THE NUMERATOR ARE
      C STORED IN R1; THE ZEROES OF THE DENOMINATOR ARE STORED IN R2.
      C THE LEADING COEFFICIENT OF BOTH POLYNOMIALS IS 1.0. THE
      C RIGHT HAND SIDE IS STORED IN V. THE SOLUTION IS COMPUTED BY
      C SOLVING ISPAN LINEAR SYSTEMS INVOLVING THE FACTORS OF THE
      C POLYNOMIAL WHICH APPEARS IN THE NUMERATOR. A CHECK FOR
      C COMMON FACTORS IS CARRIED OUT, AND THEY ARE CANCELLED IF
      C FOUND. FOUR COMPLETE LOOPS ARE PROVIDED; THE CORECT LOOP
      C FOR A GIVEN PROBLEM IS DETERMINED BY THE VALUE OF MTYPE.
      C VERSION DATE: FEBRUARY 1, 1977.
      DIMENSION TM1(1),TM2(1),V(1),U(1),W(1),E(1),R1(1),R2(1)
      COMMON /MACHEP/TOL
      TOL2=8.0E0*TOL
      MM1=M-1
      GO TO (3,50,20,75),MTYPE
      C THE LOOP FOR MTYPE = 1.
      3 DO 15 J=1,ISPAN
          D=R1(J)
          DR=R2(J)-D
          U(1)=TM1(1)+D
          Q=ABS(DR/U(1))
          IF(Q.LE.TOL2) GO TO 15
          W(1)=V(1)
          DO 5 ICOMP=2,M
              ICM1=ICOMP-1
              Q=TM2(ICOMP)/U(ICM1)
              W(ICOMP)=V(ICOMP)+W(ICM1)*Q
          5 U(ICOMP)=TM1(ICOMP)+D-Q*TM2(ICOMP)
          Q=W(M)/U(M)
          V(M)=V(M)+DR*Q
          DO 10 II=1,MM1
              ICOMP=M-II

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TRI11120
 TRI11130
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TRI20010
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 TRI20350
 TRI20360
 TRI20370

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          Q=(W(ICOMP)+TM2(ICOMP+1)*Q)/U(ICOMP)
10      V(ICOMP)=V(ICOMP)+DR*Q
15      CONTINUE
      RETURN
C      THE LOOP FOR MTYPE = 3.
20      DO 40 J=1,ISPAN
          D=R1(J)
          DR=R2(J)-D
          U(1)=TM1(1)+D
          Q=ABS(DR/U(1))
          IF(Q.LE.TOL2) GO TO 40
          W(1)=V(1)
          E(1)=TM2(1)
          DO 25 ICOMP=2,MM1
              ICM1=ICOMP-1
              Q=TM2(ICOMP)/U(ICM1)
              W(ICOMP)=V(ICOMP)+W(ICM1)*Q
              U(ICOMP)=TM1(ICOMP)+D-Q*TM2(ICOMP)
25          E(ICOMP)=Q*E(ICM1)
              E(MM1)=(E(MM1)+TM2(M))/U(MM1)
              W(MM1)=W(MM1)/U(MM1)
              DO 30 II=2,MM1
                  ICOMP=M-II
                  ICP1=ICOMP+1
                  W(ICOMP)=(W(ICOMP)+TM2(ICP1)*W(ICP1))/U(ICOMP)
30          E(ICOMP)=(E(ICOMP)+TM2(ICP1)*E(ICP1))/U(ICOMP)
              Q=TM1(M)+D-TM2(1)*E(1)-TM2(M)*E(MM1)
              Q=(V(M)+TM2(1)*W(1)+TM2(M)*W(MM1))/Q
              V(M)=V(M)+DR*Q
              DO 35 ICOMP=1,MM1
35          V(ICOMP)=V(ICOMP)+DR*(W(ICOMP)+E(ICOMP)*Q)
40      CONTINUE
      RETURN
C      THE LOOP FOR MTYPE = 2.
50      MM2=M-2
          T3=1.0E0/TM2(3)
          T2=TM2(2)*T3
          T4=TM2(M)*T3
          DO 65 J=1,ISPAN
              D=R1(J)
              DR=(R2(J)-D)*T3
              U(1)=TM2(3)/(TM1(1)+D)
              Q=ABS(DR*U(1))
              IF(Q.LE.TOL2) GO TO 65
              W(1)=V(1)
              T1=(TM1(2)+D)*T3
              Q=T2*U(1)
              W(2)=V(2)+W(1)*Q
              U(2)=1.0E0/(T1-Q*T2)
              DO 55 ICOMP=3,MM1
                  ICM1=ICOMP-1
                  W(ICOMP)=V(ICOMP)+W(ICM1)*U(ICM1)
55          U(ICOMP)=1.0E0/(T1-U(ICM1))
              Q=T4*U(MM1)
              W(M)=V(M)+W(MM1)*Q
              U(M)=(TM1(M)+D)*T3-Q*T4
              Q=W(M)/U(M)
              V(M)=V(M)+DR*Q
              Q=(W(MM1)+Q*T4)*U(MM1)
              V(MM1)=V(MM1)+DR*Q
              DO 60 II=2,MM2
                  ICOMP=M-II
                  Q=(W(ICOMP)+Q)*U(ICOMP)
60          V(ICOMP)=V(ICOMP)+DR*Q
              Q=(W(1)+T2*Q)*U(1)
              V(1)=V(1)+DR*Q
65      CONTINUE
      RETURN
C      THE LOOP FOR MTYPE = 4.
75      T2=1.0E0/TM2(2)
          M1=MM1/2
          M1P1=M1+1
          M1P2=M1+2
          M1P3=M1+3
          MSW=M-(M/2)*2
          M2=M1-MSW

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TRI20380
 TRI20390
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      RM=1.0E0
      RA=1.0E0
      IF(MSW.EQ.1) GO TO 80
      RM=2.0E0
      RA=0.0E0
80   DO 85 ICOMP=1,M1
          JCOMP=M+MSW-ICOMP
          Q=(V(JCOMP)-V(ICOMP))/2.0E0
          V(JCOMP)=(V(JCOMP)+V(ICOMP))/2.0E0
85   DO 115 J=1,ISPAN
          D=R1(J)
          DR=(R2(J)-D)*T2
          T1=(TM1(1)+D)*T2
          U(1)=1.0E0/(T1+RA)
          Q=ABS(DR*U(1))
          IF(Q.LE.TOL2) GO TO 115
          W(1)=V(1)
          DO 90 ICOMP=2,M1
              ICM1=ICOMP-1
              W(ICOMP)=V(ICOMP)+W(ICM1)*U(ICM1)
90   U(ICOMP)=1.0E0/(T1-U(ICM1))
          U(M1P1)=1.0E0/T1
          W(M1P1)=V(M1P1)
          W(M1P2)=V(M1P2)+W(M1P1)*U(M1P1)
          U(M1P2)=1.0E0/(T1-2.0E0*U(M1P1))
          DO 95 ICOMP=M1P3,MM1
              ICM1=ICOMP-1
              W(ICOMP)=V(ICOMP)+W(ICM1)*U(ICM1)
95   U(ICOMP)=1.0E0/(T1-U(ICM1))
          W(M)=V(M)+W(MM1)*U(MM1)*RM
          U(M)=T1-RM*U(MM1)-RA
          Q=W(M)/U(M)
          V(M)=V(M)+DR*Q
          DO 105 II=1,M2
              ICOMP=M-II
              Q=(W(ICOMP)+Q)*U(ICOMP)
105  V(ICOMP)=V(ICOMP)+DR*Q
          Q=(W(M1P1)+2.0E0*Q)*U(M1P1)
          V(M1P1)=V(M1P1)+DR*Q
          Q=W(M1)*U(M1)
          V(M1)=V(M1)+DR*Q
          DO 110 II=2,M1
              ICOMP=M1P1-II
              Q=(W(ICOMP)+Q)*U(ICOMP)
110  V(ICOMP)=V(ICOMP)+DR*Q
115  CONTINUE
          DO 120 ICOMP=1,M1
              JCOMP=M+MSW-ICOMP
              Q=V(JCOMP)-V(ICOMP)
              V(JCOMP)=V(JCOMP)+V(ICOMP)
120  V(ICOMP)=Q
      RETURN
      END
C
C
C   SUBROUTINE MARCH1(M,MTYPE,TM1,TM2,TN1,TN2,IBDIM,B,
1   ISPAN,LINE0,IALPHA,V1,V2)
C       IN SUBROUTINE MARCH1, THE INITIAL MARCHING PHASE OF THE
C       MARCHING ALGORITHM IS CARRIED OUT. ISPAN IS THE NUMBER OF LINES
C       PRESENT. INFORMATION ABOUT THE DIRECTION OF THE MARCH, AND
C       THE SCALARS WHICH MULTIPLY THE MARCHING VECTORS AT EACH STAGE
C       IS STORED IN IALPHA. THE BOUNDARY LINE FOR THE MARCH IS LINE0.
C       THE MARCHING VECTORS ARE V1 AND V2. THE RESULT IS RETURNED IN
C       THE VECTOR V1.
C       VERSION DATE: MARCH 1, 1977.
C       DIMENSION TM1(1),TM2(1),TN1(1),TN2(1),B(1),V1(1),V2(1)
      LINE=LINE0+IALPHA
      IBETA=(IALPHA+1)/2
      IBETA=LINE+IBETA
      BETA=TN2(IBETA)
      BETA1=-1.0E0/BETA
      LNPTR=(LINE-1)*IBDIM
      DO 5 ICOMP=1,M
          LNPTR=LNPTR+1
          V1(ICOMP)=B(LNPTR)*BETA1

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TRI21140
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5      V2(ICOMP)=0.0E0                                MCH10210
      IF(ISPAN.EQ.1) RETURN                            MCH10220
      MM1=M-1                                          MCH10230
      TC=0.0E0                                         MCH10240
      IF(MTYPE.GE.3) TC=TM2(1)                       MCH10250
      DO 20 I=2,ISPAN                                  MCH10260
          LINE=LINE+IALPHA                             MCH10270
          IBETA=IBETA+IALPHA                           MCH10280
          LNPTR=(LINE-1)*IBDIM                        MCH10290
          ALPHA=TN1(LINE)                             MCH10300
          BETA2=BETA                                   MCH10310
          BETA=TN2(IBETA)                             MCH10320
          BETA1=1.0E0/BETA                            MCH10330
          XM=V1(M)                                     MCH10340
          XB=V1(1)                                     MCH10350
          XC=XB                                        MCH10360
          TB=TC                                        MCH10370
          DO 15 ICOMP=1,MM1                            MCH10380
              LNPTR=LNPTR+1                            MCH10390
              XT=XM                                     MCH10400
              XM=XB                                     MCH10410
              XB=V1(ICOMP+1)                          MCH10420
              TT=TB                                    MCH10430
              TB=TM2(ICOMP+1)                          MCH10440
              Q=B(LNPTR)+BETA2*V2(ICOMP)+TT*XT+TB*XB  MCH10450
              V1(ICOMP)=(TM1(ICOMP)+ALPHA)*XM-Q)*BETA1 MCH10460
              V2(ICOMP)=XM                             MCH10470
15      Q=B(LNPTR+1)+BETA2*V2(M)+TB*XM+TC*XC         MCH10480
          V1(M)=(TM1(M)+ALPHA)*XB-Q)*BETA1           MCH10490
20      V2(M)=XB                                     MCH10500
      RETURN                                          MCH10510
      END                                            MCH10520

C
C
SUBROUTINE MARCH2(M,MTYPE,TM1,TM2,TN1,TN2,IBDIM,B,    MCH20010
1      ISPAN,LINE0,V1,V2)                            MCH20020
C      IN SUBROUTINE MARCH2, THE FINAL MARCHING PHASE OF THE MCH20030
C      MARCHING ALGORITHM IS CARRIED OUT. ISPAN IS THE NUMBER OF LINESMCH20040
C      PRESENT. THE BOUNDARY LINE FOR THE MARCH IS LINE0. THE MARCHINGMCH20050
C      VECTORS ARE V1 AND V2. AS THE MARCHING PROCEEDS, THE SOLUTION ISMCH20060
C      WRITTEN OVER THE RIGHT HAND SIDE B.           MCH20070
C      VERSION DATE: MARCH 1, 1977.                 MCH20080
C      DIMENSION TM1(1),TM2(1),TN1(1),TN2(1),B(1),V1(1),V2(1) MCH20090
      LINE=LINE0                                     MCH20100
      DO 5 ICOMP=1,M                                  MCH20110
          V1(ICOMP)=-V1(ICOMP)                        MCH20120
          V2(ICOMP)=0.0E0                             MCH20130
5      IF(ISPAN.EQ.1) GO TO 25                        MCH20140
      MM1=M-1                                          MCH20150
      BETA=TN2(LINE+1)                                MCH20160
      TC=0.0E0                                         MCH20170
      IF(MTYPE.GE.3) TC=TM2(1)                       MCH20180
      DO 20 I=2,ISPAN                                  MCH20190
          LINE=LINE+1                                  MCH20200
          LNPTR=(LINE-1)*IBDIM                        MCH20210
          ALPHA=TN1(LINE)                             MCH20220
          BETA2=BETA                                   MCH20230
          BETA=TN2(LINE+1)                             MCH20240
          BETA1=1.0E0/BETA                            MCH20250
          XM=V1(M)                                     MCH20260
          XB=V1(1)                                     MCH20270
          XC=XB                                        MCH20280
          TB=TC                                        MCH20290
          DO 15 ICOMP=1,MM1                            MCH20300
              LNPTR=LNPTR+1                            MCH20310
              XT=XM                                     MCH20320
              XM=XB                                     MCH20330
              XB=V1(ICOMP+1)                          MCH20340
              TT=TB                                    MCH20350
              TB=TM2(ICOMP+1)                          MCH20360
              Q=B(LNPTR)+BETA2*V2(ICOMP)+TT*XT+TB*XB  MCH20370
              V1(ICOMP)=(TM1(ICOMP)+ALPHA)*XM-Q)*BETA1 MCH20380
              V2(ICOMP)=XM                             MCH20390
15      B(LNPTR)=XM                                   MCH20400
          Q=B(LNPTR+1)+BETA2*V2(M)+TB*XM+TC*XC         MCH20410
          V1(M)=(TM1(M)+ALPHA)*XB-Q)*BETA1           MCH20420

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	V2(M)=XB	MCH20430
20	B(LNPTR+1)=XB	MCH20440
25	LNPTR=LINE*IBDIM	MCH20450
	DO 30 ICOMP=1,M	MCH20460
	LNPTR=LNPTR+1	MCH20470
30	B(LNPTR)=V1(ICOMP)	MCH20480
	RETURN	MCH20490
	END	MCH20500
C		
C		
1	SUBROUTINE MARCH3(M,MTYPE,TM1,TM2,TN1,TN2,IBDIM,B, ISPAN,LINE0,LINEC,V1,V2)	MCH30010
C	IN SUBROUTINE MARCH3, THE INITIAL MARCHING PAHSE OF THE	MCH30020
C	MARCHING ALGORITHM IS CARRIED OUT, FOR CONSTANT COEFFICIENT	MCH30030
C	PROBLEMS MEETING CERTAIN REQUIREMENTS. ISPAN IS THE NUMBER OF	MCH30040
C	LINE PRESENT. MARCHING OCCURS IN BOTH DRRECTIONS FROM THE	MCH30050
C	BOUNDARY LINE LINE0. THE MARCHING VECTORS ARE V1 AND V2.	MCH30060
C	VERSION DATE: MARCH 1, 1977.	MCH30070
C	DIMENSION TM1(1),TM2(1),TN1(1),TN2(1),B(1),V1(1),V2(1)	MCH30080
	LINE=LINE0-1	MCH30090
	LINN=LINEC+1	MCH30100
	BETA=TN2(LINE)	MCH30110
	BETA1=-1.0E0/BETA	MCH30120
	LNPTR=(LINE-1)*IBDIM	MCH30130
	NNPTR=(LINN-1)*IBDIM	MCH30140
	DO 5 ICOMP=1,M	MCH30150
	LNPTR=LNPTR+1	MCH30160
	NNPTR=NNPTR+1	MCH30170
	V1(ICOMP)=(B(LNPTR)+B(NNPTR))*BETA1	MCH30180
5	V2(ICOMP)=0.0E0	MCH30190
	IF(ISPAN.EQ.1) RETURN	MCH30200
	MM1=M-1	MCH30210
	TC=0.0E0	MCH30220
	IF(MTYPE.GE.3) TC=TM2(1)	MCH30230
	DO 20 I=2,ISPAN	MCH30240
	LINE=LINE-1	MCH30250
	LINN=LINN+1	MCH30260
	LNPTR=(LINE-1)*IBDIM	MCH30270
	NNPTR=(LINN-1)*IBDIM	MCH30280
	ALPHA=TN1(LINE)	MCH30290
	BETA2=BETA	MCH30300
	BETA=TN2(LINE)	MCH30310
	BETA1=1.0E0/BETA	MCH30320
	XM=V1(M)	MCH30330
	XB=V1(1)	MCH30340
	XC=XB	MCH30350
	TB=TC	MCH30360
	DO 15 ICOMP=1,MM1	MCH30370
	LNPTR=LNPTR+1	MCH30380
	NNPTR=NNPTR+1	MCH30390
	XT=XM	MCH30400
	XM=XB	MCH30410
	XB=V1(ICOMP+1)	MCH30420
	TT=TB	MCH30430
	TB=TM2(ICOMP+1)	MCH30440
	Q=B(LNPTR)+B(NNPTR)+BETA2*V2(ICOMP)+TT*XT+TB*XB	MCH30450
	V1(ICOMP)=((TM1(ICOMP)+ALPHA)*XM-Q)*BETA1	MCH30460
15	V2(ICOMP)=XM	MCH30470
	Q=B(LNPTR+1)+B(NNPTR+1)+BETA2*V2(M)+TB*XM+TC*XC	MCH30480
	V1(M)=((TM1(M)+ALPHA)*XB-Q)*BETA1	MCH30490
20	V2(M)=XB	MCH30500
	RETURN	MCH30510
	END	MCH30520
C		MCH30530
C		
1	SUBROUTINE GMAS(NPC,N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2, IBDIM,B,VN,VM,K,R,A)	GMAS0010
C	SUBROUTINE GMAS IS USED IN PLACE OF SUBROUTINE GMA FOR	GMAS0020
C	SOLVING, IN THE LEAST SQUARES SENSE, LINEAR SYSTEMS IN WHICH	GMAS0030
C	THE MATRIX IS POSITIVE OR NEGATIVE SEMI-DEFINITE HAVING	GMAS0040
C	ZERO AS AN EIGENVALUE OF MULTIPLICITY ONE.	GMAS0050
C	THE LINEAR SYSTEM HAS THE FORM:	GMAS0060
C		GMAS0070
C		GMAS0080
C	(TM1(I) + TN1(J)) * X(I,J)	GMAS0090
C	- TM2(I+1) * X(I+1,J) - TM2(I) * X(I-1,J)	GMAS0100
C	- TN2(J+1) * X(I,J+1) - TN2(J) * X(I,J-1) = B(I,J)	GMAS0110

C
C FOR I = 1,2,...,M, AND J = 1,2,...,N, WHERE TM2(M+1) IS GMAS0120
C INTERPRETED AS TM2(1), TN2(N+1) AS TN2(1), X(0,J) AS X(M,J), GMAS0130
C X(M+1,J) AS X(1,J), X(I,0) AS X(I,N), AND X(I,N+1) AS X(I,1). GMAS0140
C GMAS0150
C GMAS0160
C THE PARAMETER LIST: GMAS0170
C GMAS0180
C NPC IS AN INTEGER. IF NPC = 0, GMAS DOES NECESSARY GMAS0190
C PREPROCESSING CALCULATIONS. IF NPC = 1, GMAS SOLVES THE GMAS0200
C LINEAR SYSTEM. GMAS0210
C N IS AN INTEGER, AS DEFINED ABOVE. N MUST BE GREATER THAN 2. GMAS0220
C NTYPE IS AN INTEGER DESCRIBING THE ARRAYS TN1 AND TN2. GMAS0230
C NTYPE = 1: GENERAL SEPARABLE - DIRICHLET, NEUMANN, OR GMAS0240
C MIXED BOUNDARY CONDITIONS. GMAS0250
C TN1(I) IS ARBITRARY, I = 1,2,...,N. GMAS0260
C TN2(1) = 0.0; TN2(I) IS ARBITRARY, NONZERO, I = 2,3,...,N. GMAS0270
C NTYPE = 2: CONSTANT COEFFICIENT - DIRICHLET, NEUMANN, GMAS0280
C OR MIXED BOUNDARY CONDITIONS. GMAS0290
C TN1(I) = ALPHA, ALPHA AN ARBITRARY CONSTANT, I = 2,3,...,N-1. GMAS0300
C TN2(1) = 0.0; TN2(I) = BETA, BETA AN ARBITRARY NONZERO GMAS0310
C CONSTANT, I = 3,...,N-1. GMAS0320
C TOP BOUNDARY CONDITION : ONE OF GMAS0330
C TN1(1) = ALPHA; TN2(2) = BETA; (DIRICHLET) GMAS0340
C TN1(1) = ALPHA; TN2(2) = SQRT(2) * BETA; (NEUMANN-CENTERED) GMAS0350
C TN1(1) = ALPHA - BETA; TN2(2) = BETA; (NEUMANN-STAGGERED) GMAS0360
C TN1(1) = RHO, RHO ARBITRARY; TN2(2) = ZETA, ZETA ARBITRARY, GMAS0370
C NONZERO; (MIXED). GMAS0380
C BOTTOM BOUNDARY CONDITION : ONE OF GMAS0390
C TN1(N) = ALPHA; TN2(N) = BETA; (DIRICHLET) GMAS0400
C TN1(N) = ALPHA; TN2(N) = SQRT(2) * BETA; (NEUMANN-CENTERED) GMAS0410
C TN1(N) = ALPHA - BETA; TN2(N) = BETA; (NEUMANN-STAGGERED) GMAS0420
C TN1(N) = CHI, CHI ARBITRARY; TN2(N) = ETA, ETA ARBITRARY, GMAS0430
C NONZERO; (MIXED). GMAS0440
C NTYPE = 3: GENERAL SEPARABLE - PERIODIC BOUNDARY CONDITIONS. GMAS0450
C TN1(I) IS ARBITRARY, I = 1,2,...,N. GMAS0460
C TN2(I) IS ARBITRARY, NONZERO, I = 1,2,...,N. GMAS0470
C NTYPE = 4: CONSTANT COEFFICIENT - PERIODIC BOUNDARY GMAS0480
C CONDITIONS. GMAS0490
C TN1(I) = ALPHA, ALPHA ARBITRARY, I = 1,2,...,N. GMAS0500
C TN2(I) = BETA, BETA ARBITRARY, NONZERO, I = 1,2,...,N. GMAS0510
C IF N = 3, GMAS MAY RESPECIFY NTYPE. GMAS0520
C TN1 IS AN ARRAY OF LENGTH N+1, DEFINED ABOVE. GMAS0530
C TN2 IS AN ARRAY OF LENGTH N+1, DEFINED ABOVE. GMAS0540
C M IS AN INTEGER, AS DEFINED ABOVE. M MUST BE GREATER THAN 2. GMAS0550
C MTYPE IS AN INTEGER DESCRIBING THE ARRAYS TM1 AND TM2. GMAS0560
C MTYPE = 1: GENERAL SEPARABLE - DIRICHLET, NEUMANN, OR GMAS0570
C MIXED BOUNDARY CONDITIONS. GMAS0580
C MTYPE = 2: CONSTANT COEFFICIENT - DIRICHLET, NEUMANN, GMAS0590
C OR MIXED BOUNDARY CONDITIONS. GMAS0600
C MTYPE = 3: GENERAL SEPARABLE - PERIODIC BOUNDARY CONDITIONS. GMAS0610
C MTYPE = 4: CONSTANT COEFFICIENT - PERIODIC BOUNDARY GMAS0620
C CONDITIONS. GMAS0630
C RESTRICTIONS ON TM1 AND TM2 ARE ANALOGOUS TO THOSE ON GMAS0640
C TN1 AND TN2, RESPECTIVELY, FOR EACH GIVEN TYPE. GMAS0650
C IF M IS LESS THAN 6, GMAS MAY RESPECIFY MTYPE. GMAS0660
C TM1 IS AN ARRAY OF LENGTH M+1, DEFINED ABOVE. GMAS0670
C TM2 IS AN ARRAY OF LENGTH M+1, DEFINED ABOVE. GMAS0680
C IBDIM IS THE ROW DIMENSION OF THE ARRAY B, AS IT APPEARS GMAS0690
C IN THE CALLING PROGRAM. GMAS0700
C B IS A TWO DIMENSIONAL ARRAY WITH AT LEAST M ROWS AND N GMAS0710
C COLUMNS. ON INPUT, B CONTAINS THE RIGHT HAND SIDE B(I,J) GMAS0720
C AS DEFINED ABOVE. ON OUTPUT, FROM A CALL TO GMAS WITH GMAS0730
C NPC = 1, B CONTAINS THE SOLUTION X(I,J), AS DEFINED ABOVE. GMAS0740
C VN IS AN ARRAY OF LENGTH N+1. ON OUTPUT FROM A CALL TO GMAS GMAS0750
C WITH NPC = 0, VN CONTAINS THE EIGENVECTOR OF THE GMAS0760
C MATRIX TN CORRESPONDING TO THE ZERO EIGENVALUE. GMAS0770
C VM IS AN ARRAY OF LENGTH M+1. ON OUTPUT FROM A CALL TO GMAS GMAS0780
C WITH NPC = 0, VM CONTAINS THE EIGENVECTOR OF THE GMAS0790
C MATRIX TM CORRESPONDING TO THE ZERO EIGENVALUE. GMAS0800
C K IS THE MARCHING PARAMETER. MARCHING OCCURS IN THE GMAS0810
C N - DIRECTION. ON INPUT, IF K IS LESS THAN 2, IT GMAS0820
C ASSUMES THE DEFAULT VALUE K = 2. GMAS0830
C R IS AN ARRAY OF LENGTH N * P + 3 * N + 2, WHERE P IS AN GMAS0840
C INTEGER GREATER THAN OR EQUAL TO LOG2(N / K), AND K IS GMAS0850
C GREATER THAN OR EQUAL TO 2. R CONTAINS OUTPUT FROM A GMAS0860
C CALL TO GMAS WITH NPC = 0. GMAS0870

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C      A IS AN ARRAY OF LENGTH 5 * Q, WHERE Q = MAX( N+1, M+1 ).
C      A IS USED AS A SCRATCHPAD ARRAY BY GMAS.
C
C      NOTE THAT NPC, N, NTYPE, TN1, TN2, M, MTYPE, TM1, TM2 IBDIM,
C      B, K, AND R ARE IDENTICAL TO THE CORRESPONDING PARAMETERS
C      IN THE CALLING SEQUENCE FOR SUBROUTINE GMA.
C      THE PARAMETERS VN, VM, AND A ARE DIFFERENT.
C
C      GMAS CONTAINS ONE LABELED COMMON BLOCK, /MACHEP/,
C      CONTAINING ON VARIABLE, TOL. TOL IS A MACHINE DEPENDENT
C      CONSTANT EQUAL TO THE MACHINE EPSILON. IT IS INITIALIZED IN
C      GMAS IN THE FIRST EXECUTABLE STATEMENT.
C
C      A CALL TO GMAS WITH NPC = 0 REPLACES A CALL TO GMA WITH
C      NPC = 0.
C
C      A CALL TO GMAS WITH NPC = 1 REPLACES A CALL TO GMA WITH
C      WITH NPC = 1. THE RIGHT HAND SIDE B IS PROJECTED INTO THE
C      SUBSPACE ORTHOGONAL TO THE VECTOR V, WHERE V(I,J) =
C      VM(I) * VN(J). THIS IS NECESSARY TO INSURE THAT THE LINEAR
C      SYSTEM IS CONSISTENT. THE SOLUTION X IS ALSO PROJECTED INTO
C      THE SUBSPACE ORTHOGONAL TO V. ALL OTHER SOLUTIONS CAN BE
C      OBTAINED BY ADDING AN ARBITRARY SCALAR MULTIPLE OF V TO X.
C
C      ADDRESS INQUIRIES TO:
C      RANDOLPH E. BANK
C      DEPARTMENT OF MATHEMATICS
C      THE UNIVERSITY OF CHICAGO
C      CHICAGO, ILLINOIS 60637
C
C      VERSION DATE: MARCH 1, 1977.
C      DIMENSION TN1(1),TN2(1),TM1(1),TM2(1)
C      DIMENSION VN(1),VM(1),A(1),B(1),R(1)
C      COMMON /MACHEP/TOL
C      TOL=2.0E0**(-23)
C
C      CHECK N, NTYPE, M, MTYPE, AND IBDIM.
C      IF(NTYPE.GT.4).OR.(NTYPE.LT.1)) RETURN
C      IF(MTYPE.GT.4).OR.(MTYPE.LT.1)) RETURN
C      IF(N.LT.2).OR.(M.LT.2)) RETURN
C      IF(IBDIM.LT.M) RETURN
C
C      RESPECIFY NTYPE AND/OR MTYPE IF NECESSARY.
C      IF(NTYPE.EQ.2.AND.N.LE.3) NTYPE=1
C      IF(MTYPE.EQ.2.AND.M.LE.3) MTYPE=1
C      IF(NTYPE.EQ.4.AND.N.LE.3) NTYPE=3
C      IF(MTYPE.EQ.4.AND.M.LE.5) MTYPE=3
C      IF(N.LE.2) NTYPE=1
C      IF(M.LE.2) MTYPE=1
C      IF(NPC.NE.0) GO TO 15
C
C      THE PREPROCESSING CALCULATIONS.
C      KK=MIN0(K,N)
C      CALL PARTN(N,NTYPE,KK,NE,ND,R(3))
C      EPS IS ADDED TO ALL INTEGERS STORED IN THE REAL ARRAY R.
C      THIS PREVENTS UNFORTUNATE ROUNDING ERRORS FORM OCCURRING WHEN
C      THEY ARE CONVERTED BACK INTO INTEGERS.
C      EPS=0.25E0
C      R(1)=FLOAT(NE)+EPS
C      R(2)=FLOAT(ND)+EPS
C      IF(NTYPE.EQ.1.OR.NTYPE.EQ.3) CALL ROOTSG(N,NTYPE,
1      TN1,TN2,NE,ND,R(3),R(N+3),A)
C      IF(NTYPE.EQ.2.OR.NTYPE.EQ.4) CALL ROOTSC(N,NTYPE,
1      TN1,TN2,NE,ND,R(3),R(N+3),A)
C      CALL SVALUE(N,NTYPE,M,MTYPE,TM1,TM2,NE,R(N+3),TMIN)
C      J1=N+2
C      J2=J1+N+1
C      CALL SORT(N,NTYPE,NE,ND,R(3),R(N+3),A(1),A(J1),A(J2))
C      J=MAX0(N,M)+1
C      J1=J+1
C      J2=J1+J
C      J3=J2+J
C      J4=J3+J
C      IF(NTYPE.LT.3) CALL TINVIT(N,TN1,TN2,TMIN,VN,
1      A(1),A(J1),A(J2),A(J3),IFLAG)
C      IF(NTYPE.GE.3) CALL PINVIT(N,TN1,TN2,TMIN,VN,
1      A(1),A(J1),A(J2),A(J3),A(J4),IFLAG)
C      TMIN=-TMIN
C      IF(MTYPE.LT.3) CALL TINVIT(M,TM1,TM2,TMIN,VM,
GMAS0880
GMAS0890
GMAS0900
GMAS0910
GMAS0920
GMAS0930
GMAS0940
GMAS0950
GMAS0960
GMAS0970
GMAS0980
GMAS0990
GMAS1000
GMAS1010
GMAS1020
GMAS1030
GMAS1040
GMAS1050
GMAS1060
GMAS1070
GMAS1080
GMAS1090
GMAS1100
GMAS1110
GMAS1120
GMAS1130
GMAS1140
GMAS1150
GMAS1160
GMAS1170
GMAS1180
GMAS1190
GMAS1200
GMAS1210
GMAS1220
GMAS1230
GMAS1240
GMAS1250
GMAS1260
GMAS1270
GMAS1280
GMAS1290
GMAS1300
GMAS1310
GMAS1320
GMAS1330
GMAS1340
GMAS1350
GMAS1360
GMAS1370
GMAS1380
GMAS1390
GMAS1400
GMAS1410
GMAS1420
GMAS1430
GMAS1440
GMAS1450
GMAS1460
GMAS1470
GMAS1480
GMAS1490
GMAS1500
GMAS1510
GMAS1520
GMAS1530
GMAS1540
GMAS1550
GMAS1560
GMAS1570
GMAS1580
GMAS1590
GMAS1600
GMAS1610
GMAS1620
GMAS1630

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1      A(1),A(J1),A(J2),A(J3),IFLAG)
IF(MTYPE.GE.3) CALL PINVIT(M,TM1,TM2,TMIN,VM,
1      A(1),A(J1),A(J2),A(J3),A(J4),IFLAG)
RETURN
C      THE BACKSOLUTION CALCULATIONS.
15     TN2(N+1)=1.0E0
IF(NTYPE.LT.3) TN2(1)=1.0E0
NE=R(1)
ND=R(2)
CALL PROJ(1,N,NTYPE,VN,M,VM,IBDIM,B,NE,ND,R(3),COEFF)
CALL STEP1(N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2,
1      IBDIM,B,ND,R(3),R(N+3),A)
CALL PROJ(2,N,NTYPE,VN,M,VM,IBDIM,B,NE,ND,R(3),COEFF)
CALL STEP2(N,NTYPE,TN2,M,MTYPE,TM1,TM2,
1      IBDIM,B,NE,ND,R(3),R(N+3),A)
CALL PROJ(3,N,NTYPE,VN,M,VM,IBDIM,B,NE,ND,R(3),COEFF)
CALL STEP3(N,NTYPE,TN2,M,MTYPE,TM1,TM2,
1      IBDIM,B,NE,ND,R(3),R(N+3),A)
CALL STEP4(TN1,TN2,M,MTYPE,TM1,TM2,IBDIM,B,ND,R(3),R(N+3),A)
CALL PROJ(1,N,NTYPE,VN,M,VM,IBDIM,B,NE,ND,R(3),COEFF)
RETURN
END
C
C      SUBROUTINE SVALUE(N,NTYPE,M,MTYPE,TM1,TM2,NE,R,TMIN)
C      SUBROUTINE SVALUE DETERMINES TMIN, THE ROOT ASSOCIATED
C      WITH THE ZERO EIGENVALUE. THIS ROOT IS PERTURBED TO PREVENT
C      POSSIBLE DIVIDE CHECKS DURING THE BACKSOLUTION PROCESS.
C      VERSION DATE: FEBRUARY 1, 1977.
C      DIMENSION TM1(1),TM2(1),R(1)
C      COMMON /MACHEP/TOL
IF(MTYPE.LT.3) TM2(1)=0.0E0
TM2(M+1)=TM2(1)
T=0.0E0
Q=ABS(TM2(1))
DO 5 I=1,M
S=Q
Q=ABS(TM2(I+1))
T=AMAX1(T,ABS(TM1(I))+Q+S)
5      CONTINUE
L=N*NE
IF(NTYPE.GE.3) L=L+N
JMIN=L+1
IF((TM1(1)+R(JMIN)).LT.0.0E0) T=-T
L=L+N
IF(ABS(T+R(JMIN)).GT.ABS(T+R(L))) JMIN=L
TMIN=R(JMIN)
R(JMIN)=TMIN+(T+TMIN)*FLOAT(M)*TOL
RETURN
END
C
C      SUBROUTINE TINVIT(N,D,E,EIGEN,V,A,B,C,F,IFLAG)
C      SUBROUTINE TINVIT COMPUTES THE EIGENVECTOR OF A
C      SYMMETRIC TRIDIAGONAL MATRIX CORRESPONDING TO THE EIGENVALUE
C      EIGEN. INVERSE ITERATION AND GAUSSIAN ELIMINATION WITH
C      PARTIAL PIVOTING IS EMPLOYED. D AND E CONTAIN THE
C      TRIDIAGONAL MATRIX. A, B, AND C CONTAIN THE UPPER TRIANGULAR
C      MATRIX. F CONTAINS THE LOWER TRIANGULAR MATRIX.
C      V CONTAINS THE EIGENVECTOR.
C      VERSION DATE: FEBRUARY 1, 1977.
C      DIMENSION A(1),B(1),C(1),D(1),E(1),F(1),V(1)
C      COMMON /MACHEP/TOL
IFLAG=0
Q=ABS(D(1))
DO 5 I=2,N
Q=Q+ABS(D(I))+ABS(E(I))
5      EPS=TOL*Q*FLOAT(N)
Z=TOL*Q*SQRT(FLOAT(N))
E(N+1)=0.0E0
R=E(2)
U=D(1)-EIGEN
V(1)=Z
DO 15 I=2,N
IM1=I-1
V(I)=Z

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GMAS1640
GMAS1650
GMAS1660
GMAS1670
GMAS1680
GMAS1690
GMAS1700
GMAS1710
GMAS1720
GMAS1730
GMAS1740
GMAS1750
GMAS1760
GMAS1770
GMAS1780
GMAS1790
GMAS1800
GMAS1810
GMAS1820
GMAS1830
GMAS1840
GMAS1850

SVAL0010
SVAL0020
SVAL0030
SVAL0040
SVAL0050
SVAL0060
SVAL0070
SVAL0080
SVAL0090
SVAL0100
SVAL0110
SVAL0120
SVAL0130
SVAL0140
SVAL0150
SVAL0160
SVAL0170
SVAL0180
SVAL0190
SVAL0200
SVAL0210
SVAL0220
SVAL0230
SVAL0240
SVAL0250
SVAL0260

TINV0010
TINV0020
TINV0030
TINV0040
TINV0050
TINV0060
TINV0070
TINV0080
TINV0090
TINV0100
TINV0110
TINV0120
TINV0130
TINV0140
TINV0150
TINV0160
TINV0170
TINV0180
TINV0190
TINV0200
TINV0210
TINV0220
TINV0230
TINV0240

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      IF (ABS(E(I)).GT.ABS(U)) GO TO 10
      Q=E(I)/U
      F(I)=Q
      A(IM1)=U
      B(IM1)=R
      C(IM1)=0.0E0
      U=D(I)-EIGEN-Q*R
      R=E(I+1)
      GO TO 15
10    Q=U/E(I)
      F(I)=Q
      A(IM1)=-E(I)
      B(IM1)=EIGEN-D(I)
      C(IM1)=E(I+1)
      U=-R-Q*B(IM1)
      R=Q*C(IM1)
15    CONTINUE
      IF (U.EQ.0.0E0) U=EPS/FLOAT(N)
      A(N)=U
      B(N)=0.0E0
      C(N)=0.0E0
      L=1
      DO 50 J=1,5
        DO 25 I=2,N
          IM1=I-1
          Q=V(I)
          R=-E(I)
          IF (A(IM1).NE.R) GO TO 25
          Q=V(IM1)
          V(IM1)=V(I)
25    V(I)=Q+F(I)*V(IM1)
          Q=0.0E0
          R=0.0E0
          Z=0.0E0
          DO 30 II=1,N
            I=N+1-II
            V(I)=(V(I)+Q*B(I)+R*C(I))/A(I)
            R=Q
            Q=V(I)
30    Z=Z+ABS(Q)
          IF (Z.GT.1.0E0) GO TO 60
          IF (Z.NE.0.0E0) GO TO 40
          V(L)=EPS
          L=L+1
          IF (L.GT.N) L=1
          GO TO 50
40    Z=EPS/Z
          DO 45 I=1,N
            V(I)=V(I)*Z
45    CONTINUE
50    IFLAG=1
60    Z=0.0E0
          DO 65 I=1,N
            Z=Z+V(I)*V(I)
65    Z=1.0E0/SQRT(Z)
          DO 70 I=1,N
            V(I)=V(I)*Z
70    RETURN
      END
C
C
SUBROUTINE PINVIT(N,D,E,EIGEN,V,A,B,C,F,G,IFLAG)
C      SUBROUTINE PINVIT COMPUTES THE EIGENVECTOR OF A
C      SYMMETRIC MATRIX ASSOCIATED WITH PERIODIC BOUNDARY CONDITIONS,
C      CORRESPONDING TO THE EIGENVALUE EIGEN. INVERSE ITERATION
C      AND GAUSSIAN ELIMINATION WITH PARTIAL PIVOTING ARE EMPLOYED.
C      D AND E CONTAIN THE MATRIX WITH PERIODIC BOUNDARY CONDITIONS.
C      A, B, C, AND G CONTAIN THE UPPER TRIANGULAR MATRIX.
C      F CONTAINS THE LOWER TRIANGULAR MATRIX. V CONTAINS THE
C      EIGENVECTOR.
C      VERSION DATE: FEBRUARY 1, 1977.
C      DIMENSION A(1),B(1),C(1),D(1),E(1),F(1),G(1),V(1)
C      COMMON /MACHEP/TOL
      IFLAG=0
      NM1=N-1
      Q=0.0E0
      PINV0010
      PINV0020
      PINV0030
      PINV0040
      PINV0050
      PINV0060
      PINV0070
      PINV0080
      PINV0090
      PINV0100
      PINV0110
      PINV0120
      PINV0130
      PINV0140
      PINV0150

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DO 5 I=1,N
  G(I)=0.0E0
  5 Q=Q+ABS(D(I))+ABS(E(I))
  EPS=TOL*Q*FLOAT(N)
  Z=TOL*Q*SQRT(FLOAT(N))
  R=E(2)
  U=D(1)-EIGEN
  V(1)=Z
  V(N)=Z
  G(1)=E(1)
  G(NM1)=E(N)
DO 15 I=2,NM1
  IM1=I-1
  V(I)=Z
  IF(ABS(E(I)).GT.ABS(U)) GO TO 10
  Q=E(I)/U
  F(I)=Q
  A(IM1)=U
  B(IM1)=R
  C(IM1)=0.0E0
  U=D(I)-EIGEN-Q*R
  R=E(I+1)
  G(I)=G(I)+Q*G(IM1)
  GO TO 15
10 Q=U/E(I)
  F(I)=Q
  A(IM1)=-E(I)
  B(IM1)=EIGEN-D(I)
  C(IM1)=E(I+1)
  U=-R-Q*B(IM1)
  R=G(IM1)
  G(IM1)=G(I)
  G(I)=R+Q*G(IM1)
  R=Q*C(IM1)
15 CONTINUE
  IF(U.EQ.0.0E0) U=EPS/FLOAT(N)
  A(NM1)=U
  B(NM1)=0.0E0
  C(NM1)=0.0E0
  C(N-2)=0.0E0
  Q=0.0E0
  R=0.0E0
DO 20 II=1,NM1
  I=N-II
  G(I)=(G(I)+Q*B(I)+R*C(I))/A(I)
  R=Q
  Q=G(I)
20 U=D(N)-EIGEN-E(1)*G(1)-E(N)*G(NM1)
  IF(U.EQ.0.0E0) U=EPS/FLOAT(N)
  A(N)=U
  L=1
DO 50 J=1,5
  DO 25 I=2,NM1
    IM1=I-1
    Q=V(I)
    R=-E(I)
    IF(A(IM1).NE.R) GO TO 25
    Q=V(IM1)
    V(IM1)=V(I)
25 V(I)=Q+F(I)*V(IM1)
    Q=0.0E0
    R=0.0E0
DO 30 II=1,NM1
  I=N-II
  V(I)=(V(I)+Q*B(I)+R*C(I))/A(I)
  R=Q
  Q=V(I)
30 Q=(V(N)+E(1)*V(1)+E(N)*V(NM1))/A(N)
  V(N)=Q
  Z=ABS(Q)
DO 35 I=1,NM1
  V(I)=V(I)+G(I)*Q
35 Z=Z+ABS(V(I))
  IF(Z.GT.1.0E0) GO TO 60
  IF(Z.NE.0.0E0) GO TO 40
  V(L)=EPS
  PINV0160
  PINV0170
  PINV0180
  PINV0190
  PINV0200
  PINV0210
  PINV0220
  PINV0230
  PINV0240
  PINV0250
  PINV0260
  PINV0270
  PINV0280
  PINV0290
  PINV0300
  PINV0310
  PINV0320
  PINV0330
  PINV0340
  PINV0350
  PINV0360
  PINV0370
  PINV0380
  PINV0390
  PINV0400
  PINV0410
  PINV0420
  PINV0430
  PINV0440
  PINV0450
  PINV0460
  PINV0470
  PINV0480
  PINV0490
  PINV0500
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  PINV0610
  PINV0620
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  PINV0650
  PINV0660
  PINV0670
  PINV0680
  PINV0690
  PINV0700
  PINV0710
  PINV0720
  PINV0730
  PINV0740
  PINV0750
  PINV0760
  PINV0770
  PINV0780
  PINV0790
  PINV0800
  PINV0810
  PINV0820
  PINV0830
  PINV0840
  PINV0850
  PINV0860
  PINV0870
  PINV0880
  PINV0890
  PINV0900
  PINV0910

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	L=L+1	PINV0920
	IF(L.GT.N) L=1	PINV0930
	GO TO 50	PINV0940
40	Z=EPS/Z	PINV0950
	DO 45 I=1,N	PINV0960
45	V(I)=V(I)*Z	PINV0970
50	CONTINUE	PINV0980
	IFLAG=1	PINV0990
60	Z=0.0E0	PINV1000
	DO 65 I=1,N	PINV1010
65	Z=Z+V(I)*V(I)	PINV1020
	Z=1.0E0/SQRT(Z)	PINV1030
	DO 70 I=1,N	PINV1040
70	V(I)=V(I)*Z	PINV1050
	RETURN	PINV1060
	END	PINV1070
C		
C		
C	SUBROUTINE PROJ(ITYPE,N,NTYPE,VN,M,VM,IBDIM,B,NE,ND,ID,COEFF)	PROJ0010
C	SUBROUTINE PROJ COMPUTES THE PROJECTIONS NECESSARY TO	PROJ0020
C	INSURE THAT THE LINEAR SYSTEM IS CONSISTENT.	PROJ0030
C	VERSION DATE: FEBRUARY 1, 1977.	PROJ0040
	REAL ID	PROJ0050
	DIMENSION VN(1),VM(1),B(1),ID(1)	PROJ0060
	COEFF=0.0E0	PROJ0070
	GO TO (5,30,55),ITYPE	PROJ0080
5	DO 15 I=1,N	PROJ0090
	T=0.0E0	PROJ0100
	LNPTR=(I-1)*IBDIM	PROJ0110
	D	PROJ0120
	O 10 J=1,M	PROJ0130
	LNPTR=LNPTR+1	PROJ0140
10	T=T+VM(J)*B(LNPTR)	PROJ0150
15	COEFF=COEFF+VN(I)*T	PROJ0160
	DO 25 I=1,N	PROJ0170
	T=COEFF*VN(I)	PROJ0180
	LNPTR=(I-1)*IBDIM	PROJ0190
	DO 20 J=1,M	PROJ0200
	LNPTR=LNPTR+1	PROJ0210
20	B(LNPTR)=B(LNPTR)-T*VM(J)	PROJ0220
25	CONTINUE	PROJ0230
	RETURN	PROJ0240
30	NN=ND	PROJ0250
	IF(NTYPE.GE.3) NN=ND+1	PROJ0260
	ID2=ID(2)	PROJ0270
	IF((NN.LE.1).OR.(ID2.EQ.2)) RETURN	PROJ0280
	Q=0.0E0	PROJ0290
	DO 40 I=1,NN	PROJ0300
	LINE=ID(I+1)	PROJ0310
	T=0.0E0	PROJ0320
	LNPTR=(LINE-1)*IBDIM	PROJ0330
	DO 35 J=1,M	PROJ0340
	LNPTR=LNPTR+1	PROJ0350
35	T=T+VM(J)*B(LNPTR)	PROJ0360
	Q=Q+VN(LINE)*V(LINE)	PROJ0370
40	COEFF=COEFF+VN(LINE)*T	PROJ0380
	COEFF=COEFF/Q	PROJ0390
	DO 50 I=1,NN	PROJ0400
	LINE=ID(I+1)	PROJ0410
	T=COEFF*VN(LINE)	PROJ0420
	LNPTR=(LINE-1)*IBDIM	PROJ0430
	DO 45 J=1,M	PROJ0440
	LNPTR=LNPTR+1	PROJ0450
45	B(LNPTR)=B(LNPTR)-T*VM(J)	PROJ0460
50	CONTINUE	PROJ0470
	RETURN	PROJ0480
55	NN=NE-1	PROJ0490
	IF(NTYPE.GE.3) NN=NE	PROJ0500
	IF(NN.LT.0) RETURN	PROJ0510
	I=2**NN+1	PROJ0520
	I=ID(I)	PROJ0530
	LNPTR=(I-1)*IBDIM	PROJ0540
	DO 60 J=1,M	PROJ0550
	LNPTR=LNPTR+1	PROJ0560
60	COEFF=COEFF+VM(J)*B(LNPTR)	PROJ0570
	LNPTR=(I-1)*IBDIM	PROJ0580


```

DO 65 J=1,M
  LNPTR=LNPTR+1
65  B(LNPTR)=B(LNPTR)-COEFF*VM(J)
  RETURN
  END
C
C
SUBROUTINE EIGEN(N,NTYPE,D,E,K,EIG)
C   SUBROUTINE EIGEN COMPUTES THE K TH EIGENVALUE FOR
C   MATRICES OF THE TYPE ALLOWED BY SUBROUTINE GMA.
C   VERSION DATE: FEBRUARY 1, 1977.
  DIMENSION D(1),E(1)
  PI=3.1415926535897932384626433832795E0
  E(N+1)=E(N)
  IF(NTYPE.LT.3) E(1)=E(2)
  GO TO (10,15,50,60),NTYPE
10  CALL TRIEIG(N,D,E,K,EIG)
  RETURN
15  IF(N.LE.3) GO TO 10
  I=15
  T1=D(2)
  T2=E(3)
  S2=SQRT(2.0E0)*T2
  Q1=D(1)
  Q2=E(2)
  IF(Q1.EQ.T1.AND.Q2.EQ.T2) I=I-7
  IF(Q1.EQ.T1.AND.Q2.EQ.S2) I=I-6
  IF(Q1.EQ.(T1-T2).AND.Q2.EQ.T2) I=I-4
  Q1=D(N)
  Q2=E(N)
  IF(Q1.EQ.T1.AND.Q2.EQ.T2) I=I-7
  IF(Q1.EQ.T1.AND.Q2.EQ.S2) I=I-6
  IF(Q1.EQ.(T1-T2).AND.Q2.EQ.T2) I=I-4
  IF(I.GT.7) GO TO 10
  GO TO (20,25,30,35,40,10,45),I
20  EIG=T1-T2*2.0E0*COS(FLOAT(K)*PI/FLOAT(N+1))
  RETURN
25  EIG=T1-T2*2.0E0*COS(FLOAT(2*K-1)*PI/FLOAT(2*N))
  RETURN
30  EIG=T1-T2*2.0E0*COS(FLOAT(K-1)*PI/FLOAT(N-1))
  RETURN
35  EIG=T1-T2*2.0E0*COS(FLOAT(2*K-1)*PI/FLOAT(2*N+1))
  RETURN
40  EIG=T1-T2*2.0E0*COS(FLOAT(2*(K-1))*PI/FLOAT(2*N-1))
  RETURN
45  EIG=T1-T2*2.0E0*COS(FLOAT(K-1)*PI/FLOAT(N))
  RETURN
50  CALL PEREIG(N,D,E,K,EIG)
  RETURN
60  I=(K/2)*2
  EIG=D(1)-E(3)*2.0E0*COS(FLOAT(I)*PI/FLOAT(N))
  RETURN
  END
C
C
SUBROUTINE ERROR(EMARCH,BIG,TMMAX,N,TN1,TN2,K,ND)
C   SUBROUTINE ERROR PARTITIONS THE LINEAR SYSTEM USING THE
C   SAME ALGORITHM AS SUBROUTINE GMA. EMARCH IS DETERMINED BY
C   USING A ONE DIMENSIONAL MARCH WITH THE SCALAR TMMAX PLAYING THE
C   ROLE OF THE MATRIX TM.
C   VERSION DATE: FEBRUARY 1, 1977.
  DIMENSION TN1(1),TN2(1)
  EMARCH=0.0E0
  ND=N/MAX0(2,K)
  IF(ND*2.EQ.N) ND=ND-1
  NDP1=ND+1
  I1=N/NDP1
  I2=N-I1*NDP1
  I3=0
  ID=0
  DO 65 LNINDEX=1,NDP1
    LINE=ID+1
    IF(I3.LT.I2) I3=I3+1
    ID=I1*LNINDEX+I3
    IF(LNINDEX.EQ.NDP1) ID=N+1
    ISPAN=ID-1
    R1=1.0E0

```

```

PROJ0590
PROJ0600
PROJ0610
PROJ0620
EIGN0010
EIGN0020
EIGN0030
EIGN0040
EIGN0050
EIGN0060
EIGN0070
EIGN0080
EIGN0090
EIGN0100
EIGN0110
EIGN0120
EIGN0130
EIGN0140
EIGN0150
EIGN0160
EIGN0170
EIGN0180
EIGN0190
EIGN0200
EIGN0210
EIGN0220
EIGN0230
EIGN0240
EIGN0250
EIGN0260
EIGN0270
EIGN0280
EIGN0290
EIGN0300
EIGN0310
EIGN0320
EIGN0330
EIGN0340
EIGN0350
EIGN0360
EIGN0370
EIGN0380
EIGN0390
EIGN0400
EIGN0410
EIGN0420
EIGN0430
EIGN0440
EIGN0450
EIGN0460
EROR0010
EROR0020
EROR0030
EROR0040
EROR0050
EROR0060
EROR0070
EROR0080
EROR0090
EROR0100
EROR0110
EROR0120
EROR0130
EROR0140
EROR0150
EROR0160
EROR0170
EROR0180
EROR0190
EROR0200
EROR0210
EROR0220

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```

R2=0.0E0
DO 55 J=LINE,ISPAN
  IF (ABS(R1).GT.BIG) GO TO 100
  T= ((TMAX+TN1(J))*R1-TN2(J)*R2)/TN2(J+1)
  R2=R1
55  R1=T
  R1=ABS(R1)
  R2=ABS(TN2(ID)/TN2(LINE))*R1
  T=AMAX1(R1,R2)
  EMARCH=AMAX1(T,EMARCH)
65  CONTINUE
RETURN
100 EMARCH=BIG
RETURN
END
C
C
SUBROUTINE TRIEIG(N,D,E,K,EIGEN)
C   SUBROUTINE TRIEIG COMPUTES THE K TH EIGENVALUE OF AN
C   IRREDUCIBLE SYMMETRIC TRIDIAGONAL MATRIX USING BISECTION
C   AND THE STURM SEQUENCE PROPERTY. INITIAL UPPER AND LOWER
C   BOUNDS ARE DETERMINED FROM GERSCHGORIN ESTIMATES.
C   VERSION DATE: FEBRUARY 1, 1977.
  DIMENSION D(1),E(1)
  COMMON /MACHEP/TOL
  E(N+1)=0.0E0
  T=0.0E0
  XU=D(1)
  XL=D(1)
  DO 5 I=1,N
    X=T
    T=ABS(E(I+1))
    XU=AMAX1(XU,D(I)+(X+T))
    5  XL=AMIN1(XL,D(I)-(X+T))
  EPS=AMAX1(ABS(XU),ABS(XL))*TOL
  XU=XU+EPS
  XL=XL-EPS
  E(N+1)=E(N)
  E(1)=0.0E0
10  X=(XL+XU)/2.0E0
  T=EPS+4.0E0*TOL*(ABS(XL)+ABS(XU))
  IF((XU-XL).LE.T) GO TO 50
  KCOUNT=0
  T=1.0E0
  DO 20 I=1,N
    IF(T.EQ.0.0E0) T=TOL*ABS(E(I))
    T=D(I)-X-E(I)*E(I)/T
    IF(T.LT.0.0E0) KCOUNT=KCOUNT+1
20  CONTINUE
  IF(KCOUNT-K) 30,35,35
30  XL=X
  GO TO 10
35  XU=X
  GO TO 10
50  EIGEN=X
  E(1)=E(2)
  RETURN
  END
C
C
SUBROUTINE PEREIG(N,D,E,K,EIGEN)
C   SUBROUTINE PEREIG COMPUTES THE K TH EIGENVALUE OF THE
C   IRREDUCIBLE SYMMETRIC MATRIX WHICH ARISES IN THE CASE OF
C   PERIODIC BOUNDARY CONDITIONS. BISECTION AND THE STURM
C   SEQUENCE PROPERTY ARE USED. INITIAL UPPER AND LOWER BOUNDS
C   ARE DETERMINED FROM GERSCHGORIN ESTIMATES.
C   VERSION DATE: FEBRUARY 1, 1977.
  DIMENSION D(1),E(1)
  COMMON /MACHEP/TOL
  E(N+1)=E(1)
  NM1=N-1
  T=ABS(E(1))
  XU=D(1)
  XL=D(1)
  DO 5 I=1,N
    X=T
    T=ABS(E(I+1))
    XU=AMAX1(XU,D(I)+(X+T))
    5  XL=AMIN1(XL,D(I)-(X+T))
  EPS=AMAX1(ABS(XU),ABS(XL))*TOL
  XU=XU+EPS
  XL=XL-EPS
  E(N+1)=E(1)
  E(1)=0.0E0
10  X=(XL+XU)/2.0E0
  T=EPS+4.0E0*TOL*(ABS(XL)+ABS(XU))
  IF((XU-XL).LE.T) GO TO 50
  KCOUNT=0
  T=1.0E0
  DO 20 I=1,N
    IF(T.EQ.0.0E0) T=TOL*ABS(E(I))
    T=D(I)-X-E(I)*E(I)/T
    IF(T.LT.0.0E0) KCOUNT=KCOUNT+1
20  CONTINUE
  IF(KCOUNT-K) 30,35,35
30  XL=X
  GO TO 10
35  XU=X
  GO TO 10
50  EIGEN=X
  E(1)=E(2)
  RETURN
  END
C
C
SUBROUTINE TRIEIG(N,D,E,K,EIGEN)
C   SUBROUTINE TRIEIG COMPUTES THE K TH EIGENVALUE OF AN
C   IRREDUCIBLE SYMMETRIC TRIDIAGONAL MATRIX USING BISECTION
C   AND THE STURM SEQUENCE PROPERTY. INITIAL UPPER AND LOWER
C   BOUNDS ARE DETERMINED FROM GERSCHGORIN ESTIMATES.
C   VERSION DATE: FEBRUARY 1, 1977.
  DIMENSION D(1),E(1)
  COMMON /MACHEP/TOL
  E(N+1)=0.0E0
  T=0.0E0
  XU=D(1)
  XL=D(1)
  DO 5 I=1,N
    X=T
    T=ABS(E(I+1))
    XU=AMAX1(XU,D(I)+(X+T))
    5  XL=AMIN1(XL,D(I)-(X+T))
  EPS=AMAX1(ABS(XU),ABS(XL))*TOL
  XU=XU+EPS
  XL=XL-EPS
  E(N+1)=E(N)
  E(1)=0.0E0
10  X=(XL+XU)/2.0E0
  T=EPS+4.0E0*TOL*(ABS(XL)+ABS(XU))
  IF((XU-XL).LE.T) GO TO 50
  KCOUNT=0
  T=1.0E0
  DO 20 I=1,N
    IF(T.EQ.0.0E0) T=TOL*ABS(E(I))
    T=D(I)-X-E(I)*E(I)/T
    IF(T.LT.0.0E0) KCOUNT=KCOUNT+1
20  CONTINUE
  IF(KCOUNT-K) 30,35,35
30  XL=X
  GO TO 10
35  XU=X
  GO TO 10
50  EIGEN=X
  E(1)=E(2)
  RETURN
  END
C
C
SUBROUTINE PEREIG(N,D,E,K,EIGEN)
C   SUBROUTINE PEREIG COMPUTES THE K TH EIGENVALUE OF THE
C   IRREDUCIBLE SYMMETRIC MATRIX WHICH ARISES IN THE CASE OF
C   PERIODIC BOUNDARY CONDITIONS. BISECTION AND THE STURM
C   SEQUENCE PROPERTY ARE USED. INITIAL UPPER AND LOWER BOUNDS
C   ARE DETERMINED FROM GERSCHGORIN ESTIMATES.
C   VERSION DATE: FEBRUARY 1, 1977.
  DIMENSION D(1),E(1)
  COMMON /MACHEP/TOL
  E(N+1)=E(1)
  NM1=N-1
  T=ABS(E(1))
  XU=D(1)
  XL=D(1)
  DO 5 I=1,N
    X=T
    T=ABS(E(I+1))
    XU=AMAX1(XU,D(I)+(X+T))
    5  XL=AMIN1(XL,D(I)-(X+T))
  EPS=AMAX1(ABS(XU),ABS(XL))*TOL
  XU=XU+EPS
  XL=XL-EPS
  E(N+1)=E(1)
  E(1)=0.0E0
10  X=(XL+XU)/2.0E0
  T=EPS+4.0E0*TOL*(ABS(XL)+ABS(XU))
  IF((XU-XL).LE.T) GO TO 50
  KCOUNT=0
  T=1.0E0
  DO 20 I=1,N
    IF(T.EQ.0.0E0) T=TOL*ABS(E(I))
    T=D(I)-X-E(I)*E(I)/T
    IF(T.LT.0.0E0) KCOUNT=KCOUNT+1
20  CONTINUE
  IF(KCOUNT-K) 30,35,35
30  XL=X
  GO TO 10
35  XU=X
  GO TO 10
50  EIGEN=X
  E(1)=E(2)
  RETURN
  END

```

```

          T=ABS(E(I+1))
          XU=AMAX1(XU,D(I)+(X+T))
5         XL=AMINI(XL,D(I)-(X+T))
          EPS=AMAX1(ABS(XU),ABS(XL))
          EPS1=1.E-2*EPS
          EPS=EPS*TOL
          XU=XU+EPS
          XL=XL-EPS
10        X=(XL+XU)/2.0E0
          T=EPS+4.0E0*TOL*(ABS(XL)+ABS(XU))
          IF((XU-XL).LE.T) GO TO 60
          KCOUNT=0
          ITEST=0
          Q=D(1)-X
          T=E(1)
          C=D(N)-X
          IF(Q.LT.0.0E0) KCOUNT=1
          DO 30 I=2,NM1
              Q1=Q
              T1=T
              IF(Q1.EQ.0.0E0) Q1=TOL*ABS(E(I))
              V=E(I)/Q1
              Q=D(I)-X-E(I)*V
              T=T1*V
              IF(Q.LT.0.0E0) KCOUNT=KCOUNT+1
              IF(ITEST.EQ.1) GO TO 25
              IF(T1+C.EQ.C) GO TO 30
              IF(ABS(Q1).LE.EPS1) GO TO 20
              C=C-T1*T1/Q1
              GO TO 30
20          C=C-T1*T1*(D(I)-X)/(Q1*Q)
              ITEST=1
              GO TO 30
25          ITEST=0
30          CONTINUE
          IF(Q.EQ.0.0E0) Q=TOL*ABS(E(N))
          IF(ITEST) 35,35,40
35          Q=C-(E(N)+T)*((E(N)+T)/Q)
              GO TO 45
40          Q=C-E(N)*((E(N)+2.0E0*T)/Q)
45          IF(Q.LT.0.0E0) KCOUNT=KCOUNT+1
              IF(KCOUNT-K) 50,55,55
50          XL=X
              GO TO 10
55          XU=X
              GO TO 10
60          EIGEN=X
              RETURN
          END

```

C PROGRAM DRIVER
C
C THIS DRIVER GENERATES RANDOM PROBLEMS TO TEST SUBROUTINES
C KPICK, GMA, AND GMAS.
C THE USER MAY SPECIFY THE FOLLOWING PARAMETERS, WHICH APPEAR
C AS THE FIRST FIVE EXECUTABLE STATEMENTS:
C
C ITMAX = THE NUMBER OF RANDOM PROBLEMS TO BE SOLVED.
C NMAX = THE MAXIMUM VALUE OF N ALLOWED.
C MMAX = THE MAXIMUM VALUE OF M ALLOWED.
C IBDIM = THE ROW DIMENSION OF THE ARRAY B.
C TOL = THE MACHINE EPSILON.
C
C MINIMUM DIMENSIONS FOR THE ARRAYS ARE AS FOLLOWS:
C A: 5 * MAX(NMAX+1, MMAX+1)
C B: MMAX + 2 X NMAX + 2
C R: NMAX * INT(LOG2(NMAX)) + 3 * NMAX + 2
C TN1, TN2, VN: NMAX + 1
C TM1, TM2, VM: MMAX + 1
C
C THE PROGRAM IS COMPLETE EXCEPT FOR THE TIMER, WHICH IS
C REFERRED TO AS THE INTEGER FORTRAN FUNCTION ITIMER(I).
C
C VERSION DATE: MARCH 1, 1977.
C COMMON /RAN1/INIT,IINIT,ISAVE

PEIG0170
PEIG0180
PEIG0190
PEIG0200
PEIG0210
PEIG0220
PEIG0230
PEIG0240
PEIG0250
PEIG0260
PEIG0270
PEIG0280
PEIG0290
PEIG0300
PEIG0310
PEIG0320
PEIG0330
PEIG0340
PEIG0350
PEIG0360
PEIG0370
PEIG0380
PEIG0390
PEIG0400
PEIG0410
PEIG0420
PEIG0430
PEIG0440
PEIG0450
PEIG0460
PEIG0470
PEIG0480
PEIG0490
PEIG0500
PEIG0510
PEIG0520
PEIG0530
PEIG0540
PEIG0550
PEIG0560
PEIG0570
PEIG0580
PEIG0590
PEIG0600
PEIG0610
PEIG0620
PEIG0630
PEIG0640
PEIG0650

DRVR0010
DRVR0020
DRVR0030
DRVR0040
DRVR0050
DRVR0060
DRVR0070
DRVR0080
DRVR0090
DRVR0100
DRVR0110
DRVR0120
DRVR0130
DRVR0140
DRVR0150
DRVR0160
DRVR0170
DRVR0180
DRVR0190
DRVR0200
DRVR0210
DRVR0220
DRVR0230
DRVR0240
DRVR0250

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COMMON /MACHEP/TOL
DIMENSION A(160),B(33,33),R(226)
DIMENSION TN1(32),TN2(32),VN(32),TM1(32),TM2(32),VM(32)
ITMAX=100
NMAX=31
MMAX=31
IBDIM=33
TOL=2.0E0**(-23)
INIT=1367
IINIT=INIT
D1=-ALOG10(TOL)-1.0E0
DO 20 I=1,ITMAX
  N=INTRAN(3,NMAX)
  M=INTRAN(3,MMAX)
  NTYPE=INTRAN(1,4)
  MTYPE=INTRAN(1,4)
  ITYPE=0
  JTYPE=0
  IF(NTYPE.EQ.2) ITYPE=INTRAN(1,9)
  IF(MTYPE.EQ.2) JTYPE=INTRAN(1,9)
  IDEF=INTRAN(0,3)
  DEMAND=RAN(1.0E0,D1)
  CALL TEVAL(N,NTYPE,ITYPE,TN1,TN2,IDEF)
  CALL TEVAL(M,MTYPE,JTYPE,TM1,TM2,IDEF)
  ITK=ITIMER(0)
  CALL KPICK(0,N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2,
1    DEMAND,K,COND,EMARCH,DIGITS,IFLAG)
  ITK=ITIMER(0)-ITK
  IF(IDEF.GT.1) GO TO 5
  ITR=ITIMER(0)
  CALL GMA(0,N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2,IBDIM,B,
1    K,R,A)
  ITR=ITIMER(0)-ITR
  CALL RHS(N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2,IBDIM,B,
1    IDEF,VN,VM,COEFF)
  ITS=ITIMER(0)
  CALL GMA(1,N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2,IBDIM,B,
1    K,R,A)
  ITS=ITIMER(0)-ITS
  GO TO 10
5    ITR=ITIMER(0)
  CALL GMAS(0,N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2,IBDIM,B,
1    VN,VM,K,R,A)
  ITR=ITIMER(0)-ITR
  CALL RHS(N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2,IBDIM,B,
1    IDEF,VN,VM,COEFF)
  ITS=ITIMER(0)
  CALL GMAS(1,N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2,IBDIM,B,
1    VN,VM,K,R,A)
  ITS=ITIMER(0)-ITS
10   ITG=ITR+ITS
  CALL NORM(ACTUAL,N,M,IBDIM,B,IDEF,VN,VM,COEFF)
  IF(MOD(I,50).EQ.1) WRITE(6,50)
  WRITE(6,55) I,N,NTYPE,ITYPE,M,MTYPE,JTYPE,K,IFLAG,IDEF,
1    COND,EMARCH,DEMAND,DIGITS,ACTUAL,ITK,ITG,ITR,ITS
20   CONTINUE
50   FORMAT(1H1 / 46X,34HTHE GENERALIZED MARCHING ALGORITHM //
1    14X,18HPROBLEM PARAMETERS,17X,6HERRORES,13X,14HCORRECT DIGITS,
2    21X,7HTIMINGS / 4X,36HI N NTYPE M MTYPE K DEF,5X,
3    40HCOND EMARCH DEMAND DIGITS ACTUAL,4X,
3    35HKPICK GMA ROOTS SOLVE)
55   FORMAT(2X,I3,1X,I3,2X,I1,2H (,I1,1H),1X,I3,2X,I1,2H (,I1,1H),
1    2X,I2,2H (,I1,1H),2X,I2,2X,E10.3,2X,E10.3,2X,F5.2,2X,F5.2,2X,
2    F5.2,2X,I8,2X,I8,2X,I8,2X,I8)
STOP
END
C
C
FUNCTION RAN(XL,XU)
C    FUNCTION RAN(XL,XU) GENERATES PSUEDO-RANDOM REAL NUMBERS
C    ON THE CLOSED INTERVAL (XL,XU)
C    VERSION DATE: FEBRUARY 1, 1977.
COMMON /RAN1/INIT,IINIT,ISAVE
INIT=MOD(3125*INIT,65536)
DRVR0260
DRVR0270
DRVR0280
DRVR0290
DRVR0300
DRVR0310
DRVR0320
DRVR0330
DRVR0340
DRVR0350
DRVR0360
DRVR0370
DRVR0380
DRVR0390
DRVR0400
DRVR0410
DRVR0420
DRVR0430
DRVR0440
DRVR0450
DRVR0460
DRVR0470
DRVR0480
DRVR0490
DRVR0500
DRVR0510
DRVR0520
DRVR0530
DRVR0540
DRVR0550
DRVR0560
DRVR0570
DRVR0580
DRVR0590
DRVR0600
DRVR0610
DRVR0620
DRVR0630
DRVR0640
DRVR0650
DRVR0660
DRVR0670
DRVR0680
DRVR0690
DRVR0700
DRVR0710
DRVR0720
DRVR0730
DRVR0740
DRVR0750
DRVR0760
DRVR0770
DRVR0780
DRVR0790
DRVR0800
DRVR0810
DRVR0820
DRVR0830
DRVR0840
DRVR0850
DRVR0860
DRVR0870
DRVR0880
DRVR0890
DRVR0900
DRVR0910
RAN00010
RAN00020
RAN00030
RAN00040
RAN00050
RAN00060

```

	RAN=XL+(XU-XL)*FLOAT(INIT)*(2.0E0**(-16))	RAN00070
	RETURN	RAN00080
	END	RAN00090
C		
C	INTEGER FUNCTION INTRAN(IL,IU)	IRAN00100
C	FUNCTION INTRAN(IL,IU) GENERATES PSUEDO-RANDOM INTEGERS	IRAN00200
C	ON THE CLOSED INTERVAL (IL,IU)	IRAN00300
C	VERSION DATE: FEBRUARY 1, 1977.	IRAN00400
	COMMON /RAN1/IINIT,INIT,ISAVE	IRAN00500
	INIT=MOD(3125*INIT,65536)	IRAN00600
	Q=FLOAT(INIT*(IU-IL+1))*(2.0E0**(-16))	IRAN00700
	INTRAN=IL+INT(Q)	IRAN00800
	RETURN	IRAN00900
	END	IRAN01000
C		
C	SUBROUTINE TEVAL(N,NTYPE,ITYPE,D,E,IDEF)	TEVL00100
C	SUBROUTINE TEVAL GENERATES A MATRIX OF ALLOWABLE TYPE	TEVL00200
C	USING PSUEDO-RANDOM NUMBERS FOR MATRIX ELEMENTS.	TEVL00300
C	VERSION DATE: FEBRUARY 1, 1977.	TEVL00400
	DIMENSION D(1),E(1)	TEVL00500
	Q=1.0E0	TEVL00600
	IF((IDEF.EQ.1).OR.(IDEF.EQ.3)) Q=-Q	TEVL00700
	H=Q/FLOAT(N+1)	TEVL00800
	IF(IDEF.GT.1) H=0.0E0	TEVL00900
	H2=H*H*Q	TEVL01000
	X0=0.0E0	TEVL01100
	XL=Q*1.E-3	TEVL01200
	XU=Q	TEVL01300
	NP1=N+1	TEVL01400
	IF(NTYPE.EQ.2.OR.NTYPE.EQ.4) GO TO 15	TEVL01500
	DO 5 I=1,NP1	TEVL01600
5	E(I)=RAN(XL,XU)	TEVL01700
	IF(NTYPE.EQ.3) E(NP1)=E(1)	TEVL01800
	DO 10 I=1,N	TEVL01900
10	D(I)=E(I)+E(I+1)+RAN(X0,H2)	TEVL02000
	IF((IDEF.LE.1).OR.(NTYPE.EQ.3)) RETURN	TEVL02100
	D(1)=E(2)	TEVL02200
	D(N)=E(N)	TEVL02300
	RETURN	TEVL02400
15	T1=RAN(XL,XU)	TEVL02500
	T2=2.0E0*T1+RAN(X0,H2)	TEVL02600
	DO 20 I=1,NP1	TEVL02700
	D(I)=T2	TEVL02800
20	E(I)=T1	TEVL02900
	IF(NTYPE.EQ.4) RETURN	TEVL03000
	S2=SQRT(2.0E0)*T1	TEVL03100
	IF(IDEF.GT.1) ITYPE=5	TEVL03200
	IF(ITYPE.LE.3) GO TO 30	TEVL03300
	IF(ITYPE.GT.6) GO TO 25	TEVL03400
	L=INTRAN(0,1)	TEVL03500
	IF(L.EQ.0) E(2)=S2	TEVL03600
	IF(L.EQ.1) D(1)=T2-T1	TEVL03700
	GO TO 30	TEVL03800
25	E(2)=S2	TEVL03900
	D(1)=T2+RAN(X0,H)	TEVL04000
30	I=ITYPE-(ITYPE/3)*3	TEVL04100
	IF(I.EQ.1) RETURN	TEVL04200
	IF(I.EQ.0) GO TO 40	TEVL04300
	L=INTRAN(0,1)	TEVL04400
	IF(L.EQ.0) E(N)=S2	TEVL04500
	IF(L.EQ.1) D(N)=T2-T1	TEVL04600
	RETURN	TEVL04700
40	E(N)=S2	TEVL04800
	D(N)=T2+RAN(X0,H)	TEVL04900
	RETURN	TEVL05000
	END	TEVL05100
C		TEVL05200
C		
C		
1	SUBROUTINE RHS(N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2,IBDIM,B,	RHS00010
	IDEF,VN,VM,COEFF)	RHS00020
C	SUBROUTINE RHS GENERATES A RIGHT HAND SIDE FOR SUBROUTINE	RHS00030
C	GMA USING PSUEDO-RANDOM NUMBERS. ISAVE STORES THE CURRENT VALUE	RHS00040


```

15 DO 20 I=1,M                                NORM0240
      T=VM(I)*COEFF                            NORM0250
DO 20 J=1,N                                NORM0260
      U=RAN(XL,XU)-T*VN(J)                    NORM0270
      S=DBLE(B(I,J))-DBLE(U)                  NORM0280
      E1=E1+S*S                                NORM0290
20      E2=E2+U*U                              NORM0300
25 E1=E1/E2                                    NORM0310
      IF(E1.LT.1.0E0) DIGITS=-ALOG10(E1)/2.0E0 NORM0320
      RETURN                                    NORM0330
      END                                        NORM0340

C
C
      SUBROUTINE KPICK(NPC,N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2,    KPIK0010
1      DEMAND,K,COND,EMARCH,DIGITS,IFLAG)                   KPIK0020
C      SUBROUTINE KPICK IS DESIGNED TO AID THE USER IN SELECTING KPIK0030
C      THE MARCHING PARAMETER K TO BE USED AS INPUT FOR SUBROUTINES KPIK0040
C      GMA AND GMAS.                                        KPIK0050
C                                                         KPIK0060
C      THE PARAMETER LIST:                                KPIK0070
C                                                         KPIK0080
C      NPC IS AN INTEGER. IF NPC = 0, SUBROUTINE KPICK DETERMINES K KPIK0090
C      SUCH THAT THE SOLUTIONS COMPUTED USING SUBROUTINE GMA (GMAS) KPIK0100
C      WILL HAVE APPROXIMATELY 'DEMAND' SIGNIFICANT DIGITS. KPIK0110
C      IF NPC = 1, SUBROUTINE KPICK WILL TEST THE INPUT VALUE KPIK0120
C      OF K.                                              KPIK0130
C      N,NTYPE,TN1,TN2,M,MTYPE,TM1,TM2 ARE IDENTICAL TO THE KPIK0140
C      CORRESPONDING PARAMETERS IN THE CALLING SEQUENCE FOR KPIK0150
C      SUBROUTINE GMA (GMAS).                             KPIK0160
C      DEMAND IS A REAL NUMBER, STATING THE NUMBER OF SIGNIFICANT KPIK0170
C      DIGITS DESIRED IN THE COMPUTED SOLUTION. IT MUST BE KPIK0180
C      SPECIFIED ON INPUT IF NPC = 0.                    KPIK0190
C      K IS AN INTEGER. IF NPC = 0, ON OUTPUT K IS EQUAL TO THE KPIK0200
C      MARCHING PARAMETER DETERMINED BY KPICK. IF NPC = 1, THE KPIK0210
C      USER MUST SPECIFY THE VALUE OF K TO BE TESTED AS AN KPIK0220
C      INPUT PARAMETER.                                  KPIK0230
C      COND IS A REAL NUMBER, NORMALLY RETURNING THE CONDITION KPIK0240
C      NUMBER OF THE LINEAR SYSTEM.                       KPIK0250
C      EMARCH IS A REAL NUMBER, NORMALLY RETURNING AN ESTIMATE OF KPIK0260
C      THE ERROR DUE TO MARCHING FOR THE OUTPUT VALUE OF THE KPIK0270
C      MARCHING PARAMETER K.                             KPIK0280
C      DIGITS IS A REAL NUMBER, NORMALLY RETURNING A VALUE OF KPIK0290
C      MAX(0.0, -ALOG10( ( COND + EMARCH ) * TOL ), WHERE TOL KPIK0300
C      IS EQUAL TO THE MACHINE EPSILON. THIS IS TYPICALLY THE KPIK0310
C      NUMBER OF SIGNIFICANT DIGITS (IN THE 2 - NORM) ONE MAY KPIK0320
C      EXPECT IN THE COMPUTED SOLUTION FOR THE GIVEN VALUE OF K. KPIK0330
C      IFLAG IS AN INTEGER DESCRIBING ERROR RETURNS.     KPIK0340
C      IFLAG = 0: NORMAL RETURN.                          KPIK0350
C      IFLAG = 1: KPICK FAILED TO SUCCESSFULLY COMPUTE COND. THE KPIK0360
C      ALGORITHM USED BY KPICK ASSUMES THAT THE LINEAR KPIK0370
C      SYSTEM IS EITHER POSITIVE OR NEGATIVE DEFINITE. KPIK0380
C      IF THE LINEAR SYSTEM IS POSITIVE OR NEGATIVE KPIK0390
C      SEMI-DEFINITE WITH A ZERO EIGENVALUE OF MULTIPLICITY KPIK0400
C      ONE, KPICK COMPUTES THE CONDITION NUMBER RELATIVE TO KPIK0410
C      THE SUBSPACE ORTHOGONAL TO THE EIGENVECTOR ASSOCIATED KPIK0420
C      WITH THE ZERO EIGENVALUE. OTHERWISE, KPICK RETURNS KPIK0430
C      THE DEFAULT VALUES COND = EMARCH = DIGITS = 0.0, KPIK0440
C      AND K = 2 IF NPC = 0.                               KPIK0450
C      IFLAG = 2: KPICK HAS DETERMINED THAT THE MARCHING ERROR MAY KPIK0460
C      NOT SATISFY THE ASSUMPTIONS UNDERLYING THE ALGORITHM KPIK0470
C      USED TO COMPUTE EMARCH, AND THUS THE COMPUTED VALUE KPIK0480
C      OF DIGITS MAY BE IN DOUBT. IF NPC = 0, KPICK RETURNS KPIK0490
C      ESTIMATES FOR K = 2. IF NPC = 1, KPICK RETURNS KPIK0500
C      ESTIMATES FOR THE INPUT VALUE OF K.               KPIK0510
C      IFLAG = 3: CONDITIONS 1 AND 2 EXIST.               KPIK0520
C      IFLAG = 4: THE REQUESTED DEMAND CANNOT BE SATISFIED BY ANY KPIK0530
C      VALUE OF K GREATER THAN OR EQUAL TO 2. KPICK RETURNS KPIK0540
C      ESTIMATES FOR K = 2. THIS ERROR RETURN CAN OCCUR ONLY KPIK0550
C      IF NPC = 0.                                        KPIK0560
C      IFLAG = 5: CONDITIONS 1 AND 4 EXIST.               KPIK0570
C      IFLAG = 6: CONDITIONS 2 AND 4 EXIST.               KPIK0580
C      IFLAG = 7: CONDITIONS 1, 2, AND 4 EXIST.           KPIK0590
C      IFLAG = 8: THE PARAMETERS N, NTYPE, TN1, AND/OR TN2 ARE KPIK0600
C      INCORRECTLY SPECIFIED.                             KPIK0610
C      IFLAG = 9: THE PARAMETERS M, MTYPE, TM1, AND/OR TM2 ARE KPIK0620
C      INCORRECTLY SPECIFIED.                             KPIK0630

```

```

C          IFLAG = 10: CONDITIONS 8 AND 9 EXIST.
C
C          IF IFLAG = 0, 2, 4, 6, THEN THE LINEAR SYSTEM IS SUITABLE
C          FOR SUBROUTINE GMA, WITH THE QUALIFICATIONS NOTED ABOVE.
C          IF IFLAG = 1, 3, 5, 7, WITH OTHER THAN DEFAULT VALUES FOR
C          COND, EMARCH, AND DIGITS, THEN THE LINEAR SYSTEM IS SUITABLE
C          FOR SUBROUTINE GMAS, WITH THE QUALIFICATIONS NOTED ABOVE.
C
C          KPICK CONTAINS ONE LABELED COMMON BLOCK, /MACHEP/,
C          CONTAINING ON VARIABLE, TOL. TOL IS A MACHINE DEPENDENT
C          CONSTANT EQUAL TO THE MACHINE EPSILON. IT IS INITIALIZED IN
C          KPICK IN THE FIRST EXECUTABLE STATEMENT.
C
C          ADDRESS INQUIRIES TO:
C          RANDOLPH E. BANK
C          DEPARTMENT OF MATHEMATICS
C          THE UNIVERSITY OF CHICAGO
C          CHICAGO, ILLINOIS 60637
C
C          VERSION DATE: FEBRUARY 1, 1977.
C          DIMENSION TN1(1),TN2(1),TM1(1),TM2(1)
C          COMMON /MACHEP/TOL
C          TOL=2.0E0**(-23)
C          ESTABLISH DEFAULT VALUES
C          IFLAG=0
C          IF(NPC.NE.1) K=2
C          COND=0.0E0
C          EMARCH=0.0E0
C          DIGITS=0.0E0
C          NN=N
C          IF(NTYPE.GE.3) NN=N-1
C          CHECK MATRIX SPECIFICATIONS.
C          JFLAG=8
C          CALL TCHECK(N,NTYPE,TN1,TN2,IFLAG,JFLAG)
C          JFLAG=9+IFLAG/8
C          CALL TCHECK(M,MTYPE,TM1,TM2,IFLAG,JFLAG)
C          IF(IFLAG.NE.0) RETURN
C          DETERMINE IF THE SYSTEM IS POSITIVE OR NEGATIVE DEFINITE.
C          NS=2
C          MS=2
C          CALL EIGEN(N,NTYPE,TN1,TN2,1,TNMIN)
C          CALL EIGEN(N,NTYPE,TN1,TN2,N,TNMAX)
C          CALL EIGEN(M,MTYPE,TM1,TM2,1,TMMIN)
C          CALL EIGEN(M,MTYPE,TM1,TM2,M,TMMAX)
C          IF(ABS(TMMAX+TNMAX).GT.ABS(TMMIN+TNMAX)) GO TO 5
C          MS=M-1
C          T1=TMMAX
C          TMMAX=TMMIN
C          TMMIN=T1
C          5 IF(ABS(TMMAX+TNMAX).GT.ABS(TMMAX+TNMIN)) GO TO 10
C          NS=N-1
C          T1=TNMAX
C          TNMAX=TNMIN
C          TNMIN=T1
C          10 T1=TMMAX+TNMAX
C          T2=TMMIN+TNMIN
C          IF(T2/T1.GT.TOL) GO TO 15
C          IF THE MATRIX IS NOT POSITIVE OR NEGATIVE DEFINITE,
C          DETERMINE IF IT IS POSITIVE OR NEGATIVE SEMI-DEFINITE
C          WITH A ZERO EIGENVALUE OF MULTIPLICITY ONE.
C          IFLAG=1
C          IF(ABS(T2).GT.ABS(T1)*TOL) RETURN
C          CALL EIGEN(N,NTYPE,TN1,TN2,NS,TNSNG)
C          CALL EIGEN(M,MTYPE,TM1,TM2,MS,TMSNG)
C          T2=TMMIN+TNSNG
C          T3=TMSNG+TNMIN
C          IF(ABS(T2).GT.ABS(T3)) T2=T3
C          IF(T2/T1.LE.TOL) RETURN
C          DETERMINE THE CONDITION NUMBER OF THE LINEAR SYSTEM.
C          THE VALUE GF + SQRT( GF * GF - 1. ) IS AN ESTIMATE OF THE
C          EXPONENTIAL GROWTH PER MARCHING STEP.
C          15 COND=T1/T2
C          GF=(2.0E0*TMMAX+TNMAX+TNMIN)/(TNMAX-TNMIN)
C          BIG=COND/TOL
C          IF(NPC.EQ.1) GO TO 50
C          ESTIMATE AN UPPER BOUND FOR K, KMAX, BASED ON AN EXPONENTIAL

```

```

KPIK0640
KPIK0650
KPIK0660
KPIK0670
KPIK0680
KPIK0690
KPIK0700
KPIK0710
KPIK0720
KPIK0730
KPIK0740
KPIK0750
KPIK0760
KPIK0770
KPIK0780
KPIK0790
KPIK0800
KPIK0810
KPIK0820
KPIK0830
KPIK0840
KPIK0850
KPIK0860
KPIK0870
KPIK0880
KPIK0890
KPIK0900
KPIK0910
KPIK0920
KPIK0930
KPIK0940
KPIK0950
KPIK0960
KPIK0970
KPIK0980
KPIK0990
KPIK1000
KPIK1010
KPIK1020
KPIK1030
KPIK1040
KPIK1050
KPIK1060
KPIK1070
KPIK1080
KPIK1090
KPIK1100
KPIK1110
KPIK1120
KPIK1130
KPIK1140
KPIK1150
KPIK1160
KPIK1170
KPIK1180
KPIK1190
KPIK1200
KPIK1210
KPIK1220
KPIK1230
KPIK1240
KPIK1250
KPIK1260
KPIK1270
KPIK1280
KPIK1290
KPIK1300
KPIK1310
KPIK1320
KPIK1330
KPIK1340
KPIK1350
KPIK1360
KPIK1370
KPIK1380
KPIK1390

```


COLLECTED ALGORITHMS (cont.)**527-P44- 0**

```
                IF (ABS(D1-D(I)).GT.D2) GO TO 50
                IF (ABS(E1-E(I)).GT.E2) GO TO 50
40  CONTINUE
    IF (E(2).EQ.0.0E0) GO TO 50
    IF (E(N).EQ.0.0E0) GO TO 50
    RETURN
50  IFLAG=JFLAG
    RETURN
    END
```

TCHK0310
TCHK0320
TCHK0330
TCHK0340
TCHK0350
TCHK0360
TCHK0370
TCHK0380
TCHK0390

ALGORITHM 528

Framework for a Portable Library [Z]

P. A. FOX, A. D. HALL, and N. L. SCHRYER

Bell Laboratories

Key Words and Phrases: portability, mathematical libraries, error handling, storage management, memory allocation, machine dependencies

CR Categories: 4.4, 5.1

Language: Fortran

DESCRIPTION

The three program packages presented here provide a framework for a portable Fortran subroutine library. They were developed for the Bell Laboratories library PORT [1]. The packages are: machine-dependent constants, automatic error handling, and dynamic storage allocation using a stack. There are interdependencies among the packages in the sense that the error handling is used by both the others, and it in turn uses machine-dependent constants provided by the first. However, care is taken to avoid any actual or apparent recursion.

Two non-ANSI Standard Fortran assumptions are made in the algorithm. The first is that there is no runtime subscript range checking; the second is that variables (local to a subprogram) initialized by DATA statements, and then changed within the subprogram, keep their values from one invocation of the subprogram to the next. Appendix A of [1] discusses nonstandard usage in more detail.

REFERENCES

1. FOX, P.A., HALL, A.D., AND SCHRYER, N.L. The PORT mathematical subroutine library. *ACM Trans. Math. Software* 4, 2 (June 1978), 104-126.

Machine-Dependent Constants

The first package contains three Fortran function subprograms which can be invoked to determine basic machine or operating system dependent constants. Values are provided in commented DATA statements for the Burroughs 5700/6700/7700, the CDC 6000/7000 series, the Data General Eclipse, DEC PDP 10 (KA and KI processors), the DEC PDP 11, the Harris S220, the Honeywell 6000 series, the IBM 360/370 series, the SEL systems 85/86, the Univac 1100 series, the XEROX SIGMA 5/7/9; others can be added. When the library is moved to a new environment, only the appropriate DATA statements in these three subprograms need to be activated by removing the C's from column 1.

The three functions are: IIMACH, which delivers integer constants, RIMACH,

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which delivers single-precision floating-point (REAL) constants, and DIMACH, which delivers double-precision floating-point constants. These functions have a single integer argument indicating the particular constant desired. For example, RIMACH(2) is the largest single-precision floating-point number on the host machine, so the statement

```
XMAX = RIMACH(2)
```

sets XMAX to this largest number.

The machine-dependent values provided are defined in the listing given here. (Please note that the systems for which the machine constants are given are printed here in alphabetical order but that their corresponding sequence numbers are not in numerical order.) Details of constant specification and usage are described in [1].

ALGORITHM

[Only the PORT utility programs are included here. A Users' Guide and test programs are available upon request from the ACM Algorithms Distribution Service.]

```

      INTEGER FUNCTION IIMACH(I)
C
C   I/O UNIT NUMBERS.
C
C   IIMACH( 1) = THE STANDARD INPUT UNIT.
C
C   IIMACH( 2) = THE STANDARD OUTPUT UNIT.
C
C   IIMACH( 3) = THE STANDARD PUNCH UNIT.
C
C   IIMACH( 4) = THE STANDARD ERROR MESSAGE UNIT.
C
C   WORDS.
C
C   IIMACH( 5) = THE NUMBER OF BITS PER INTEGER STORAGE UNIT.
C
C   IIMACH( 6) = THE NUMBER OF CHARACTERS PER INTEGER STORAGE UNIT.
C
C   INTEGERS.
C
C   ASSUME INTEGERS ARE REPRESENTED IN THE S-DIGIT, BASE-A FORM
C
C           SIGN ( X(S-1)*A**(S-1) + ... + X(1)*A + X(0) )
C
C           WHERE 0 .LE. X(I) .LT. A FOR I=0,...,S-1.
C
C   IIMACH( 7) = A, THE BASE.
C
C   IIMACH( 8) = S, THE NUMBER OF BASE-A DIGITS.
C
C   IIMACH( 9) = A**S - 1, THE LARGEST MAGNITUDE.
C
C   FLOATING-POINT NUMBERS.
C
C   ASSUME FLOATING-POINT NUMBERS ARE REPRESENTED IN THE T-DIGIT,
C   BASE-B FORM
C
C           SIGN (B**E)*( (X(1)/B) + ... + (X(T)/E**T) )
C
C           WHERE 0 .LE. X(I) .LT. B FOR I=1,...,T,
C           0 .LT. X(1), AND EMIN .LE. E .LE. EMAX.
C
C   IIMACH(10) = B, THE BASE.
C
C   SINGLE-PRECISION
C
C   IIMACH(11) = T, THE NUMBER OF BASE-B DIGITS.
C
      MCHI0000
      MCHI0020
      MCHI0040
      MCHI0060
      MCHI0080
      MCHI0100
      MCHI0120
      MCHI0140
      MCHI0160
      MCHI0180
      MCHI0200
      MCHI0220
      MCHI0240
      MCHI0260
      MCHI0280
      MCHI0300
      MCHI0320
      MCHI0340
      MCHI0360
      MCHI0380
      MCHI0400
      MCHI0420
      MCHI0440
      MCHI0460
      MCHI0480
      MCHI0500
      MCHI0520
      MCHI0540
      MCHI0560
      MCHI0580
      MCHI0600
      MCHI0620
      MCHI0640
      MCHI0660
      MCHI0680
      MCHI0700
      MCHI0720
      MCHI0740
      MCHI0760
      MCHI0780
      MCHI0800
      MCHI0820
      MCHI0840
      MCHI0860
      MCHI0880
      MCHI0900
      MCHI0920
      MCHI0940

```

```

C      IIMACH(12) = EMIN, THE SMALLEST EXPONENT E.          MCHI0960
C                                                                MCHI0980
C      IIMACH(13) = EMAX, THE LARGEST EXPONENT E.          MCHI1000
C                                                                MCHI1020
C      DOUBLE-PRECISION                                     MCHI1040
C                                                                NCHI1060
C      IIMACH(14) = T, THE NUMBER OF BASE-B DIGITS.        MCHI1080
C                                                                MCHI1100
C      IIMACH(15) = EMIN, THE SMALLEST EXPONENT E.        MCHI1120
C                                                                MCHI1140
C      IIMACH(16) = EMAX, THE LARGEST EXPONENT E.          MCHI1160
C                                                                MCHI1180
C      TO ALTER THIS FUNCTION FOR A PARTICULAR ENVIRONMENT, MCHI1200
C      THE DESIRED SET OF DATA STATEMENTS SHOULD BE ACTIVATED BY MCHI1220
C      REMOVING THE C FROM COLUMN 1. ALSO, THE VALUES OF MCHI1240
C      IIMACH(1) - IIMACH(4) SHOULD BE CHECKED FOR CONSISTENCY MCHI1260
C      WITH THE LOCAL OPERATING SYSTEM.                    MCHI1280
C                                                                MCHI1300
C      INTEGER IMACH(16),OUTPUT                             MCHI1320
C                                                                MCHI1340
C      EQUIVALENCE (IMACH(4),OUTPUT)                       MCHI1360
C                                                                MCHI1380
C      MACHINE CONSTANTS FOR THE BURROUGHS 1/00 SYSTEM.    MCHI4880
C                                                                MCHI4900
C      DATA IMACH( 1) / 7 /                                MCHI4920
C      DATA IMACH( 2) / 2 /                                MCHI4940
C      DATA IMACH( 3) / 2 /                                MCHI4960
C      DATA IMACH( 4) / 2 /                                MCHI4980
C      DATA IMACH( 5) / 36 /                               MCHI5000
C      DATA IMACH( 6) / 4 /                                MCHI5020
C      DATA IMACH( 7) / 2 /                                MCHI5040
C      DATA IMACH( 8) / 33 /                               MCHI5060
C      DATA IMACH( 9) / Z1FFFFFFF /                       MCHI5080
C      DATA IMACH(10) / 2 /                                MCHI5100
C      DATA IMACH(11) / 24 /                               MCHI5120
C      DATA IMACH(12) / -256 /                             MCHI5140
C      DATA IMACH(13) / 255 /                             MCHI5160
C      DATA IMACH(14) / 60 /                               MCHI5180
C      DATA IMACH(15) / -256 /                             MCHI5200
C      DATA IMACH(16) / 255 /                             MCHI5220
C                                                                MCHI5240
C      MACHINE CONSTANTS FOR THE BURROUGHS 5700 SYSTEM.    MCHI4500
C                                                                MCHI4520
C      DATA IMACH( 1) / 5 /                                MCHI4540
C      DATA IMACH( 2) / 6 /                                MCHI4560
C      DATA IMACH( 3) / 7 /                                MCHI4580
C      DATA IMACH( 4) / 6 /                                MCHI4600
C      DATA IMACH( 5) / 48 /                               MCHI4620
C      DATA IMACH( 6) / 6 /                                MCHI4640
C      DATA IMACH( 7) / 2 /                                MCHI4660
C      DATA IMACH( 8) / 39 /                               MCHI4680
C      DATA IMACH( 9) / 0000777777777777 /                MCHI4700
C      DATA IMACH(10) / 8 /                                MCHI4720
C      DATA IMACH(11) / 13 /                               MCHI4740
C      DATA IMACH(12) / -50 /                              MCHI4760
C      DATA IMACH(13) / 76 /                               MCHI4780
C      DATA IMACH(14) / 26 /                               MCHI4800
C      DATA IMACH(15) / -50 /                              MCHI4820
C      DATA IMACH(16) / 76 /                               MCHI4840
C                                                                MCHI4860
C      MACHINE CONSTANTS FOR THE BURROUGHS 6700/7700 SYSTEMS. MCHI4120
C                                                                MCHI4140
C      DATA IMACH( 1) / 5 /                                MCHI4160
C      DATA IMACH( 2) / 6 /                                MCHI4180
C      DATA IMACH( 3) / 7 /                                MCHI4200
C      DATA IMACH( 4) / 6 /                                MCHI4220
C      DATA IMACH( 5) / 48 /                               MCHI4240
C      DATA IMACH( 6) / 6 /                                MCHI4260
C      DATA IMACH( 7) / 2 /                                MCHI4280
C      DATA IMACH( 8) / 39 /                               MCHI4300
C      DATA IMACH( 9) / 0000777777777777 /                MCHI4320
C      DATA IMACH(10) / 8 /                                MCHI4340
C      DATA IMACH(11) / 13 /                               MCHI4360
C      DATA IMACH(12) / -50 /                              MCHI4380
C      DATA IMACH(13) / 76 /                               MCHI4400
C      DATA IMACH(14) / 26 /                               MCHI4420

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C	DATA IMACH(15) / -32754 /	MCHI4440
C	DATA IMACH(16) / 32780 /	MCHI4460
C		MCHI4480
C	MACHINE CONSTANTS FOR THE CDC 6000/7000 SERIES.	MCHI2180
C		MCHI2200
C	DATA IMACH(1) / 5 /	MCHI2220
C	DATA IMACH(2) / 6 /	MCHI2240
C	DATA IMACH(3) / 7 /	MCHI2260
C	DATA IMACH(4) / 6 /	MCHI2280
C	DATA IMACH(5) / 60 /	MCHI2300
C	DATA IMACH(6) / 10 /	MCHI2320
C	DATA IMACH(7) / 2 /	MCHI2340
C	DATA IMACH(8) / 48 /	MCHI2360
C	DATA IMACH(9) / 00007777777777777777B /	MCHI2380
C	DATA IMACH(10) / 2 /	MCHI2400
C	DATA IMACH(11) / 48 /	MCHI2420
C	DATA IMACH(12) / -974 /	MCHI2440
C	DATA IMACH(13) / 1070 /	MCHI2460
C	DATA IMACH(14) / 96 /	MCHI2480
C	DATA IMACH(15) / -927 /	MCHI2500
C	DATA IMACH(16) / 1070 /	MCHI2520
C		MCHI2540
C	MACHINE CONSTANTS FOR THE CRAY 1	
C		
C	DATA IMACH(1) / 100 /	
C	DATA IMACH(2) / 101 /	
C	DATA IMACH(3) / 102 /	
C	DATA IMACH(4) / 101 /	
C	DATA IMACH(5) / 64 /	
C	DATA IMACH(6) / 8 /	
C	DATA IMACH(7) / 2 /	
C	DATA IMACH(8) / 63 /	
C	DATA IMACH(9) / 7777777777777777777777777777B /	
C	DATA IMACH(10) / 2 /	
C	DATA IMACH(11) / 48 /	
C	DATA IMACH(12) / -8192 /	
C	DATA IMACH(13) / 8191 /	
C	DATA IMACH(14) / 96 /	
C	DATA IMACH(15) / -8192 /	
C	DATA IMACH(16) / 8191 /	
C		
C	MACHINE CONSTANTS FOR THE DATA GENERAL ECLIPSE S/200	MCHI5720
C		MCHI5740
C	DATA IMACH(1) / 11 /	MCHI5760
C	DATA IMACH(2) / 12 /	MCHI5780
C	DATA IMACH(3) / 8 /	MCHI5800
C	DATA IMACH(4) / 10 /	MCHI5820
C	DATA IMACH(5) / 16 /	MCHI5840
C	DATA IMACH(6) / 2 /	MCHI5860
C	DATA IMACH(7) / 2 /	MCHI5880
C	DATA IMACH(8) / 15 /	MCHI5900
C	DATA IMACH(9) / 32767 /	MCHI5920
C	DATA IMACH(10) / 16 /	MCHI5940
C	DATA IMACH(11) / 6 /	MCHI5960
C	DATA IMACH(12) / -64 /	MCHI5980
C	DATA IMACH(13) / 63 /	MCHI6000
C	DATA IMACH(14) / 14 /	MCHI6020
C	DATA IMACH(15) / -64 /	MCHI6040
C	DATA IMACH(16) / 63 /	MCHI6060
C		MCHI6080
C	MACHINE CONSTANTS FOR THE HARRIS 220	MCHI6100
C		MCHI6120
C	DATA IMACH(1) / 5 /	MCHI6140
C	DATA IMACH(2) / 6 /	MCHI6160
C	DATA IMACH(3) / 0 /	MCHI6180
C	DATA IMACH(4) / 6 /	MCHI6200
C	DATA IMACH(5) / 24 /	MCHI6220
C	DATA IMACH(6) / 3 /	MCHI6240
C	DATA IMACH(7) / 2 /	MCHI6260
C	DATA IMACH(8) / 23 /	MCHI6280
C	DATA IMACH(9) / 8388607 /	MCHI6300
C	DATA IMACH(10) / 2 /	MCHI6320
C	DATA IMACH(11) / 23 /	MCHI6340
C	DATA IMACH(12) / -127 /	MCHI6360
C	DATA IMACH(13) / 127 /	MCHI6380
C	DATA IMACH(14) / 38 /	MCHI6400

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C DATA IMACH(15) / -127 / MCHI6420
C DATA IMACH(16) / 127 / MCHI6440
C MCHI6460
C MACHINE CONSTANTS FOR THE HONEYWELL 6000/6000 SERIES. MCHI1400
C MCHI1420
C DATA IMACH( 1) / 5 / MCHI1440
C DATA IMACH( 2) / 6 / MCHI1460
C DATA IMACH( 3) / 43 / MCHI1480
C DATA IMACH( 4) / 6 / MCHI1500
C DATA IMACH( 5) / 36 / MCHI1520
C DATA IMACH( 6) / 6 / MCHI1540
C DATA IMACH( 7) / 2 / MCHI1560
C DATA IMACH( 8) / 35 / MCHI1580
C DATA IMACH( 9) / 037777777777 / MCHI1600
C DATA IMACH(10) / 2 / MCHI1620
C DATA IMACH(11) / 27 / MCHI1640
C DATA IMACH(12) / -127 / MCHI1660
short course: intro to tso, part 2, thursday 16 march, 7-9 pm, room 354 ml OPER
C DATA IMACH(13) / 127 / MCHI1680
C DATA IMACH(14) / 63 / MCHI1700
C DATA IMACH(15) / -127 / MCHI1720
C DATA IMACH(16) / 127 / MCHI1740
C MCHI1760
C MACHINE CONSTANTS FOR THE IBM 360/370 SERIES, MCHI1780
C THE XEROX SIGMA 5/7/9 AND THE SEL SYSTEMS 85/86. MCHI1800
C MCHI1820
C DATA IMACH( 1) / 5 / MCHI1840
C DATA IMACH( 2) / 6 / MCHI1860
C DATA IMACH( 3) / 7 / MCHI1880
C DATA IMACH( 4) / 6 / MCHI1900
C DATA IMACH( 5) / 32 / MCHI1920
C DATA IMACH( 6) / 4 / MCHI1940
C DATA IMACH( 7) / 2 / MCHI1960
C DATA IMACH( 8) / 31 / MCHI1980
C DATA IMACH( 9) / Z7FFFFFFF / MCHI2000
C DATA IMACH(10) / 16 / MCHI2020
C DATA IMACH(11) / 6 / MCHI2040
C DATA IMACH(12) / -64 / MCHI2060
C DATA IMACH(13) / 63 / MCHI2080
C DATA IMACH(14) / 14 / MCHI2100
C DATA IMACH(15) / -64 / MCHI2120
C DATA IMACH(16) / 63 / MCHI2140
C MCHI2160
C MACHINE CONSTANTS FOR THE PDP-10 (KA PROCESSOR). MCHI2180
C MCHI2200
C DATA IMACH( 1) / 5 / MCHI2220
C DATA IMACH( 2) / 6 / MCHI2240
C DATA IMACH( 3) / 5 / MCHI2260
C DATA IMACH( 4) / 6 / MCHI2280
C DATA IMACH( 5) / 36 / MCHI2300
C DATA IMACH( 6) / 5 / MCHI2320
C DATA IMACH( 7) / 2 / MCHI2340
C DATA IMACH( 8) / 35 / MCHI2360
C DATA IMACH( 9) / "377777777777 / MCHI2380
C DATA IMACH(10) / 2 / MCHI2400
C DATA IMACH(11) / 27 / MCHI2420
C DATA IMACH(12) / -128 / MCHI2440
C DATA IMACH(13) / 127 / MCHI2460
C DATA IMACH(14) / 54 / MCHI2480
C DATA IMACH(15) / -101 / MCHI2500
C DATA IMACH(16) / 127 / MCHI2520
C MCHI2540
C MACHINE CONSTANTS FOR THE PDP-10 (KI PROCESSOR). MCHI2560
C MCHI2580
C DATA IMACH( 1) / 5 / MCHI2600
C DATA IMACH( 2) / 6 / MCHI2620
C DATA IMACH( 3) / 5 / MCHI2640
C DATA IMACH( 4) / 6 / MCHI2660
C DATA IMACH( 5) / 36 / MCHI2680
C DATA IMACH( 6) / 5 / MCHI2700
C DATA IMACH( 7) / 2 / MCHI2720
C DATA IMACH( 8) / 35 / MCHI2740
C DATA IMACH( 9) / "377777777777 / MCHI2760
C DATA IMACH(10) / 2 / MCHI2780
C DATA IMACH(11) / 27 / MCHI2800
C DATA IMACH(12) / -128 / MCHI2820
C DATA IMACH(13) / 127 / MCHI2840
C DATA IMACH(14) / 54 / MCHI2860
C DATA IMACH(15) / -101 / MCHI2880
C DATA IMACH(16) / 127 / MCHI2900
C MCHI2920
C MACHINE CONSTANTS FOR THE PDP-10 (KI PROCESSOR). MCHI2940
C MCHI2960
C DATA IMACH( 1) / 5 / MCHI2980
C DATA IMACH( 2) / 6 / MCHI3000
C DATA IMACH( 3) / 5 / MCHI3020
C DATA IMACH( 4) / 6 / MCHI3040
C DATA IMACH( 5) / 36 / MCHI3060
C DATA IMACH( 6) / 5 / MCHI3080
C DATA IMACH( 7) / 2 / MCHI3100
C DATA IMACH( 8) / 35 / MCHI3120
C DATA IMACH( 9) / "377777777777 / MCHI3140
C DATA IMACH(10) / 2 / MCHI3160
C DATA IMACH(11) / 27 / MCHI3180
C DATA IMACH(12) / -128 / MCHI3200

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C DATA IMACH(13) / 127 / MCHI3220
C DATA IMACH(14) / 62 / MCHI3240
C DATA IMACH(15) / -128 / MCHI3260
C DATA IMACH(16) / 127 / MCHI3280
C MCHI3300
C MACHINE CONSTANTS FOR PDP-11 FORTRAN'S SUPPORTING MCHI3320
C 32-BIT INTEGER ARITHMETIC. MCHI3340
C MCHI3360
C DATA IMACH( 1) / 5 / MCHI3380
C DATA IMACH( 2) / 6 / MCHI3400
C DATA IMACH( 3) / 5 / MCHI3420
C DATA IMACH( 4) / 6 / MCHI3440
C DATA IMACH( 5) / 32 / MCHI3460
C DATA IMACH( 6) / 4 / MCHI3480
C DATA IMACH( 7) / 2 / MCHI3500
C DATA IMACH( 8) / 31 / MCHI3520
C DATA IMACH( 9) / 2147483647 / MCHI3540
C DATA IMACH(10) / 2 / MCHI3560
C DATA IMACH(11) / 24 / MCHI3580
C DATA IMACH(12) / -127 / MCHI3600
C DATA IMACH(13) / 127 / MCHI3620
C DATA IMACH(14) / 56 / MCHI3640
C DATA IMACH(15) / -127 / MCHI3660
C DATA IMACH(16) / 127 / MCHI3680
C MCHI3700
C MACHINE CONSTANTS FOR PDP-11 FORTRAN'S SUPPORTING MCHI3720
C 16-BIT INTEGER ARITHMETIC. MCHI3740
C MCHI3760
C DATA IMACH( 1) / 5 / MCHI3780
C DATA IMACH( 2) / 6 / MCHI3800
C DATA IMACH( 3) / 5 / MCHI3820
C DATA IMACH( 4) / 6 / MCHI3840
C DATA IMACH( 5) / 16 / MCHI3860
C DATA IMACH( 6) / 2 / MCHI3880
C DATA IMACH( 7) / 2 / MCHI3900
C DATA IMACH( 8) / 15 / MCHI3920
C DATA IMACH( 9) / 32767 / MCHI3940
C DATA IMACH(10) / 2 / MCHI3960
C DATA IMACH(11) / 24 / MCHI3980
C DATA IMACH(12) / -127 / MCHI4000
C DATA IMACH(13) / 127 / MCHI4020
C DATA IMACH(14) / 56 / MCHI4040
C DATA IMACH(15) / -127 / MCHI4060
C DATA IMACH(16) / 127 / MCHI4080
C MCHI4100
C MACHINE CONSTANTS FOR THE UNIVAC 1100 SERIES. MCHI5260
C MCHI5280
C NOTE THAT THE PUNCH UNIT, IIMACH(3), HAS BEEN SET TO 7 MCHI5300
C WHICH IS APPROPRIATE FOR THE UNIVAC-FOR SYSTEM. MCHI5320
C IF YOU HAVE THE UNIVAC-FTN SYSTEM, SET IT TO 1. MCHI5340
C MCHI5360
C DATA IMACH( 1) / 5 / MCHI5380
C DATA IMACH( 2) / 6 / MCHI5400
C DATA IMACH( 3) / 7 / MCHI5420
C DATA IMACH( 4) / 6 / MCHI5440
C DATA IMACH( 5) / 36 / MCHI5460
C DATA IMACH( 6) / 6 / MCHI5480
C DATA IMACH( 7) / 2 / MCHI5500
C DATA IMACH( 8) / 35 / MCHI5520
C DATA IMACH( 9) / 037777777777 / MCHI5540
C DATA IMACH(10) / 2 / MCHI5560
C DATA IMACH(11) / 27 / MCHI5580
C DATA IMACH(12) / -128 / MCHI5600
C DATA IMACH(13) / 127 / MCHI5620
C DATA IMACH(14) / 60 / MCHI5640
C DATA IMACH(15) / -1024 / MCHI5660
C DATA IMACH(16) / 1023 / MCHI5680
C MCHI5700
C IF (I .LT. 1 .OR. I .GT. 16) GO TO 10 MCHI6480
C MCHI6500
C IIMACH=IMACH(I) MCHI6520
C RETURN MCHI6540
C MCHI6560
C 10 WRITE(OUTPUT,90000) MCHI6580
C 90000 FORMAT(39H1ERROR 1 IN IIMACH - I OUT OF BOUNDS) MCHI6600
C MCHI6620

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```

      CALL FDUMP
C
      STOP
C
      END
MCHI6640
MCHI6660
MCHI6680
MCHI6700
MCHI6720

```

The function for REAL floating-point numbers is

```

      REAL FUNCTION RIMACH(I)
C
C SINGLE-PRECISION MACHINE CONSTANTS
C
C RIMACH(1) = B**(EMIN-1), THE SMALLEST POSITIVE MAGNITUDE.
C
C RIMACH(2) = B**EMAX*(1 - B**(-T)), THE LARGEST MAGNITUDE.
C
C RIMACH(3) = B**(-T), THE SMALLEST RELATIVE SPACING.
C
C RIMACH(4) = B**(1-T), THE LARGEST RELATIVE SPACING.
C
C RIMACH(5) = LOG10(B)
C
C TO ALTER THIS FUNCTION FOR A PARTICULAR ENVIRONMENT,
C THE DESIRED SET OF DATA STATEMENTS SHOULD BE ACTIVATED BY
C REMOVING THE C FROM COLUMN 1.
C
C WHERE POSSIBLE, OCTAL OR HEXADEcimal CONSTANTS HAVE BEEN USED
C TO SPECIFY THE CONSTANTS EXACTLY WHICH HAS IN SOME CASES
C REQUIRED THE USE OF EQUIVALENT INTEGER ARRAYS.
C
      INTEGER SMALL(2)
      INTEGER LARGE(2)
      INTEGER RIGHT(2)
      INTEGER DIVER(2)
      INTEGER LOG10(2)
C
      REAL RMACH(5)
C
      EQUIVALENCE (RMACH(1),SMALL(1))
      EQUIVALENCE (RMACH(2),LARGE(1))
      EQUIVALENCE (RMACH(3),RIGHT(1))
      EQUIVALENCE (RMACH(4),DIVER(1))
      EQUIVALENCE (RMACH(5),LOG10(1))
C
      MACHINE CONSTANTS FOR THE BURROUGHS 1700 SYSTEM.
C
      DATA RMACH(1) / Z400800000 /
      DATA RMACH(2) / Z5FFFFFFF /
      DATA RMACH(3) / Z4E9800000 /
      DATA RMACH(4) / Z4EA800000 /
      DATA RMACH(5) / Z500E730E8 /
C
      MACHINE CONSTANTS FOR THE BURROUGHS 5700/6700/7700 SYSTEMS.
C
      DATA RMACH(1) / 017710000000000000 /
      DATA RMACH(2) / 0077777777777777 /
      DATA RMACH(3) / 013110000000000000 /
      DATA RMACH(4) / 013010000000000000 /
      DATA RMACH(5) / 01157163034761675 /
C
      MACHINE CONSTANTS FOR THE CDC 6000/7000 SERIES.
C
      DATA RMACH(1) / 00014000000000000000B /
      DATA RMACH(2) / 377677777777777777B /
      DATA RMACH(3) / 16404000000000000000B /
      DATA RMACH(4) / 16414000000000000000B /
      DATA RMACH(5) / 17164642023241175720B /
C
      MACHINE CONSTANTS FOR THE CRAY 1
C
      DATA RMACH(1) / 2000040000000000000000B /
      DATA RMACH(2) / 57777777777777777777B /
      DATA RMACH(3) / 3772140000000000000000B /
MCHR0000
MCHR0020
MCHR0040
MCHR0060
MCHR0080
MCHR0100
MCHR0120
MCHR0140
MCHR0160
MCHR0180
MCHR0200
MCHR0220
MCHR0240
MCHR0260
MCHR0280
MCHR0300
MCHR0320
MCHR0340
MCHR0360
MCHR0380
MCHR0400
MCHR0420
MCHR0440
MCHR0460
MCHR0480
MCHR0500
MCHR0520
MCHR0540
MCHR0560
MCHR0580
MCHR0600
MCHR0620
MCHR0640
MCHR0660
MCHR0680
MCHR0700
MCHR2140
MCHR2160
MCHR2180
MCHR2200
MCHR2220
MCHR2240
MCHR2260
MCHR2280
MCHR1980
MCHR2000
MCHR2020
MCHR2040
MCHR2060
MCHR2080
MCHR2100
MCHR2120
MCHR1060
MCHR1080
MCHR1100
MCHR1120
MCHR1140
MCHR1160
MCHR1180
MCHR1200

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C   DATA RMACH(4) / 37722400000000000000B /
C   DATA RMACH(5) / 377774642023241175720B /
C
C   MACHINE CONSTANTS FOR THE DATA GENERAL ECLIPSE S/200
C
C   NOTE - IT MAY BE APPROPRIATE TO INCLUDE THE FOLLOWING CARD -
C   STATIC RMACH(5)
C
C   DATA SMALL/20K,0/,LARGE/77777K,177777K/
C   DATA RIGHT/35420K,0/,DIVER/36020K,0/
C   DATA LOG10/40423K,42023K/
C
C   MACHINE CONSTANTS FOR THE HARRIS 220
C
C   DATA SMALL(1),SMALL(2) / '20000000, '00000201 /
C   DATA LARGE(1),LARGE(2) / '37777777, '00000177 /
C   DATA RIGHT(1),RIGHT(2) / '20000000, '00000352 /
C   DATA DIVER(1),DIVER(2) / '20000000, '00000353 /
C   DATA LOG10(1),LOG10(2) / '23210115, '00000377 /
C
C   MACHINE CONSTANTS FOR THE HONEYWELL 600/6000 SERIES.
C
C   DATA RMACH(1) / 0402400000000 /
C   DATA RMACH(2) / 037677777777 /
C   DATA RMACH(3) / 0714400000000 /
C   DATA RMACH(4) / 0716400000000 /
C   DATA RMACH(5) / 0776464202324 /
C
C   MACHINE CONSTANTS FOR THE IBM 360/370 SERIES,
C   THE XEROX SIGMA 5/7/9 AND THE SEL SYSTEMS 85/86.
C
C   DATA RMACH(1) / Z00100000 /
C   DATA RMACH(2) / Z7FFFFFFF /
C   DATA RMACH(3) / Z3B100000 /
C   DATA RMACH(4) / Z3C100000 /
C   DATA RMACH(5) / Z41134413 /
C
C   MACHINE CONSTANTS FOR THE PDP-10 (KA OR KI PROCESSOR).
C
C   DATA RMACH(1) / "000400000000 /
C   DATA RMACH(2) / "37777777777 /
C   DATA RMACH(3) / "146400000000 /
C   DATA RMACH(4) / "147400000000 /
C   DATA RMACH(5) / "177464202324 /
C
C   MACHINE CONSTANTS FOR PDP-11 FORTRAN'S SUPPORTING
C   32-BIT INTEGERS (EXPRESSED IN INTEGER AND OCTAL).
C
C   DATA SMALL(1) / 8388608 /
C   DATA LARGE(1) / 2147483647 /
C   DATA RIGHT(1) / 880803840 /
C   DATA DIVER(1) / 889192448 /
C   DATA LOG10(1) / 1067065499 /
C
C   DATA RMACH(1) / 000040000000 /
C   DATA RMACH(2) / 01777777777 /
C   DATA RMACH(3) / 006440000000 /
C   DATA RMACH(4) / 006500000000 /
C   DATA RMACH(5) / 007746420233 /
C
C   MACHINE CONSTANTS FOR PDP-11 FORTRAN'S SUPPORTING
C   16-BIT INTEGERS (EXPRESSED IN INTEGER AND OCTAL).
C
C   DATA SMALL(1),SMALL(2) / 128, 0 /
C   DATA LARGE(1),LARGE(2) / 32767, -1 /
C   DATA RIGHT(1),RIGHT(2) / 13440, 0 /
C   DATA DIVER(1),DIVER(2) / 13568, 0 /
C   DATA LOG10(1),LOG10(2) / 16282, 8347 /
C
C   DATA SMALL(1),SMALL(2) / 0000200, 0000000 /
C   DATA LARGE(1),LARGE(2) / 0077777, 0177777 /
C   DATA RIGHT(1),RIGHT(2) / 0032200, 0000000 /
C   DATA DIVER(1),DIVER(2) / 0032400, 0000000 /
C   DATA LOG10(1),LOG10(2) / 0037632, 0020233 /
C
MCHR2460
MCHR2480
MCHR2500
MCHR2520
MCHR2540
MCHR2560
MCHR2580
MCHR2600
MCHR2620
MCHR2640
MCHR2660
MCHR2680
MCHR2700
MCHR2720
MCHR2740
MCHR2760
MCHR2780
MCHR0720
MCHR0740
MCHR0760
MCHR0780
MCHR0800
MCHR0820
MCHR0840
MCHR0860
MCHR0880
MCHR0900
MCHR0920
MCHR0940
MCHR0960
MCHR0980
MCHR1000
MCHR1020
MCHR1040
MCHR1220
MCHR1240
MCHR1260
MCHR1280
MCHR1300
MCHR1320
MCHR1340
MCHR1360
MCHR1380
MCHR1400
MCHR1420
MCHR1440
MCHR1460
MCHR1480
MCHR1500
MCHR1520
MCHR1540
MCHR1560
MCHR1580
MCHR1600
MCHR1620
MCHR1640
MCHR1660
MCHR1680
MCHR1700
MCHR1720
MCHR1740
MCHR1760
MCHR1780
MCHR1800
MCHR1820
MCHR1840
MCHR1860
MCHR1880
MCHR1900
MCHR1920
MCHR1940
MCHR1960

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```

C     MACHINE CONSTANTS FOR THE UNIVAC 1100 SERIES.
C
C     DATA RMACH(1) / 00004000000000 /
C     DATA RMACH(2) / 037777777777 /
C     DATA RMACH(3) / 01464000000000 /
C     DATA RMACH(4) / 01474000000000 /
C     DATA RMACH(5) / 0177464202324 /
C
C     IF (I .LT. 1 .OR. I .GT. 5)
1     CALL SETERR(24HRIMACH - I OUT OF BOUNDS,24,1,2)
C
C     RIMACH = RMACH(I)
C     RETURN
C
C     END

```

```

MCHR2300
MCHR2320
MCHR2340
MCHR2360
MCHR2380
MCHR2400
MCHR2420
MCHR2440
MCHR2800
MCHR2820
MCHR2840
MCHR2860
MCHR2880
MCHR2900
MCHR2920

```

The function for DOUBLE-PRECISION floating-point numbers is

```

C     DOUBLE PRECISION FUNCTION DIMACH(I)
C
C     DOUBLE-PRECISION MACHINE CONSTANTS
C
C     DIMACH( 1) = B**(EMIN-1), THE SMALLEST POSITIVE MAGNITUDE.
C
C     DIMACH( 2) = B**EMAX*(1 - B**(-T)), THE LARGEST MAGNITUDE.
C
C     DIMACH( 3) = B**(-T), THE SMALLEST RELATIVE SPACING.
C
C     DIMACH( 4) = B**(1-T), THE LARGEST RELATIVE SPACING.
C
C     DIMACH( 5) = LOG10(B)
C
C     TO ALTER THIS FUNCTION FOR A PARTICULAR ENVIRONMENT,
C     THE DESIRED SET OF DATA STATEMENTS SHOULD BE ACTIVATED BY
C     REMOVING THE C FROM COLUMN 1.
C
C     WHERE POSSIBLE, OCTAL OR HEXADECIMAL CONSTANTS HAVE BEEN USED
C     TO SPECIFY THE CONSTANTS EXACTLY WHICH HAS IN SOME CASES
C     REQUIRED THE USE OF EQUIVALENT INTEGER ARRAYS.
C
C     INTEGER SMALL(4)
C     INTEGER LARGE(4)
C     INTEGER RIGHT(4)
C     INTEGER DIVER(4)
C     INTEGER LOG10(4)
C
C     DOUBLE PRECISION DMACH(5)
C
C     EQUIVALENCE (DMACH(1),SMALL(1))
C     EQUIVALENCE (DMACH(2),LARGE(1))
C     EQUIVALENCE (DMACH(3),RIGHT(1))
C     EQUIVALENCE (DMACH(4),DIVER(1))
C     EQUIVALENCE (DMACH(5),LOG10(1))
C
C     MACHINE CONSTANTS FOR THE BURROUGHS 1700 SYSTEM.
C
C     DATA SMALL(1) / ZC00800000 /
C     DATA SMALL(2) / Z000000000 /
C
C     DATA LARGE(1) / ZDFFFFFFF /
C     DATA LARGE(2) / ZFFFFFFF /
C
C     DATA RIGHT(1) / ZCC5800000 /
C     DATA RIGHT(2) / Z000000000 /
C
C     DATA DIVER(1) / ZCC6800000 /
C     DATA DIVER(2) / Z000000000 /
C
C     DATA LOG10(1) / ZD00E730E7 /
C     DATA LOG10(2) / ZC77800DC0 /
C
C     MACHINE CONSTANTS FOR THE BURROUGHS 5700 SYSTEM.
C
C     DATA SMALL(1) / 0177100000000000 /
C     DATA SMALL(2) / 0000000000000000 /
C
C

```

```

MCHD0000
MCHD0020
MCHD0040
MCHD0060
MCHD0080
MCHD0100
MCHD0120
MCHD0140
MCHD0160
MCHD0180
MCHD0200
MCHD0220
MCHD0240
MCHD0260
MCHD0280
MCHD0300
MCHD0320
MCHD0340
MCHD0360
MCHD0380
MCHD0400
MCHD0420
MCHD0440
MCHD0460
MCHD0480
MCHD0500
MCHD0520
MCHD0540
MCHD0560
MCHD0580
MCHD0600
MCHD0620
MCHD0640
MCHD0660
MCHD0680
MCHD0700
MCHD3360
MCHD3380
MCHD3400
MCHD3420
MCHD3440
MCHD3460
MCHD3480
MCHD3500
MCHD3520
MCHD3540
MCHD3560
MCHD3580
MCHD3600
MCHD3620
MCHD3640
MCHD3660
MCHD3680
MCHD3020
MCHD3040
MCHD3060
MCHD3080
MCHD3100

```

C	DATA LARGE(1) / 0077777777777777 /	MCHD3120
C	DATA LARGE(2) / 0000777777777777 /	MCHD3140
C		MCHD3160
C	DATA RIGHT(1) / 0146100000000000 /	MCHD3180
C	DATA RIGHT(2) / 0000000000000000 /	MCHD3200
C		MCHD3220
C	DATA DIVER(1) / 0145100000000000 /	MCHD3240
C	DATA DIVER(2) / 0000000000000000 /	MCHD3260
C		MCHD3280
C	DATA LOG10(1) / 01157163034761674 /	MCHD3300
C	DATA LOG10(2) / 00006677466732724 /	MCHD3320
C		MCHD3340
C	MACHINE CONSTANTS FOR THE BURROUGHS 6700/7700 SYSTEMS.	MCHD2680
C		MCHD2700
C	DATA SMALL(1) / 0177100000000000 /	MCHD2720
C	DATA SMALL(2) / 0777000000000000 /	MCHD2740
C		MCHD2760
C	DATA LARGE(1) / 0077777777777777 /	MCHD2780
C	DATA LARGE(2) / 0777777777777777 /	MCHD2800
C		MCHD2820
C	DATA RIGHT(1) / 0146100000000000 /	MCHD2840
C	DATA RIGHT(2) / 0000000000000000 /	MCHD2860
C		MCHD2880
C	DATA DIVER(1) / 0145100000000000 /	MCHD2900
C	DATA DIVER(2) / 0000000000000000 /	MCHD2920
C		MCHD2940
C	DATA LOG10(1) / 01157163034761674 /	MCHD2960
C	DATA LOG10(2) / 00006677466732724 /	MCHD2980
C		MCHD3000
C	MACHINE CONSTANTS FOR THE CDC 6000/7000 SERIES.	MCHD1060
C		MCHD1080
C	DATA SMALL(1) / 006040000000000000B /	MCHD1100
C	DATA SMALL(2) / 000000000000000000B /	MCHD1120
C		MCHD1140
C	DATA LARGE(1) / 377677777777777777B /	MCHD1160
C	DATA LARGE(2) / 371677777777777777B /	MCHD1180
C		MCHD1200
C	DATA RIGHT(1) / 156040000000000000B /	MCHD1220
C	DATA RIGHT(2) / 150000000000000000B /	MCHD1240
C		MCHD1260
C	DATA DIVER(1) / 156140000000000000B /	MCHD1280
C	DATA DIVER(2) / 150100000000000000B /	MCHD1300
C		MCHD1320
C	DATA LOG10(1) / 17164642023241175717B /	MCHD1340
C	DATA LOG10(2) / 16367571421742254654B /	MCHD1360
C		MCHD1380
C	MACHINE CONSTANTS FOR THE CRAY 1	
C		
C	DATA SMALL(1) / 20000400000000000000B /	
C	DATA SMALL(2) / 00000000000000000000B /	
C		
C	DATA LARGE(1) / 57777777777777777777B /	
C	DATA LARGE(2) / 0000077777777777777777B /	
C	DATA RIGHT(1) / 3772140000000000000000B /	
C	DATA RIGHT(2) / 0000000000000000000000B /	
C		
C	DATA DIVER(1) / 3772240000000000000000B /	
C	DATA DIVER(2) / 0000000000000000000000B /	
C		
C	DATA LOG10(1) / 377774642023241175717B /	
C	DATA LOG10(2) / 000007571421742254654B /	
C		
C	MACHINE CONSTANTS FOR THE DATA GENERAL ECLIPSE S/200	MCHD3880
C		MCHD3900
C	NOTE - IT MAY BE APPROPRIATE TO INCLUDE THE FOLLOWING CARD -	MCHD3920
C	STATIC DMACH(5)	MCHD3940
C		MCHD3960
C	DATA SMALL/20K,3*0/,LARGE/77777K,3*17777K/	MCHD3980
C	DATA RIGHT/31420K,3*0/,DIVER/32020K,3*0/	MCHD4000
C	DATA LOG10/40423K,42023K,50237K,74776K/	MCHD4020
C		MCHD4040
C	MACHINE CONSTANTS FOR THE HARRIS 220	MCHD4060
C		MCHD4080
C	DATA SMALL(1),SMALL(2) / '20000000, '00000201 /	MCHD4100
C	DATA LARGE(1),LARGE(2) / '37777777, '3777577 /	MCHD4120

```

C DATA RIGHT(1),RIGHT(2) / '20000000, '00000333 / MCHD4140
C DATA DIVER(1),DIVER(2) / '20000000, '00000334 / MCHD4160
C DATA LOG10(1),LOG10(2) / '23210115, '10237777 / MCHD4180
C
C MACHINE CONSTANTS FOR THE HONEYWELL 600/6000 SERIES. MCHD0720
C MCHD0740
C DATA SMALL(1),SMALL(2) / 0402400000000, 0000000000000 / MCHD0760
C DATA LARGE(1),LARGE(2) / 0376777777777, 0777777777777 / MCHD0780
C DATA RIGHT(1),RIGHT(2) / 0604400000000, 0000000000000 / MCHD0800
C DATA DIVER(1),DIVER(2) / 0606400000000, 0000000000000 / MCHD0820
C DATA LOG10(1),LOG10(2) / 0776464202324, 0117571775714 / MCHD0840
C MCHD0860
C MACHINE CONSTANTS FOR THE IBM 360/370 SERIES, MCHD0880
C THE XEROX SIGMA 5/7/9 AND THE SEL SYSTEMS 85/86. MCHD0900
C MCHD0920
C DATA SMALL(1),SMALL(2) / Z00100000, Z00000000 / MCHD0940
C DATA LARGE(1),LARGE(2) / Z7FFFFFFF, ZFFFFFFF / MCHD0960
C DATA RIGHT(1),RIGHT(2) / Z33100000, Z00000000 / MCHD0980
C DATA DIVER(1),DIVER(2) / Z34100000, Z00000000 / MCHD1000
C DATA LOG10(1),LOG10(2) / Z41134413, Z509F79FF / MCHD1020
C MCHD1040
C MACHINE CONSTANTS FOR THE PDP-10 (KA PROCESSOR). MCHD1400
C MCHD1420
C DATA SMALL(1),SMALL(2) / "033400000000, "000000000000 / MCHD1440
C DATA LARGE(1),LARGE(2) / "377777777777, "344777777777 / MCHD1460
C DATA RIGHT(1),RIGHT(2) / "113400000000, "000000000000 / MCHD1480
C DATA DIVER(1),DIVER(2) / "114400000000, "000000000000 / MCHD1500
C DATA LOG10(1),LOG10(2) / "177464202324, "144117571776 / MCHD1520
C MCHD1540
C MACHINE CONSTANTS FOR THE PDP-10 (KI PROCESSOR). MCHD1560
C MCHD1580
C DATA SMALL(1),SMALL(2) / "000400000000, "000000000000 / MCHD1600
C DATA LARGE(1),LARGE(2) / "377777777777, "377777777777 / MCHD1620
C DATA RIGHT(1),RIGHT(2) / "103400000000, "000000000000 / MCHD1640
C DATA DIVER(1),DIVER(2) / "104400000000, "000000000000 / MCHD1660
C DATA LOG10(1),LOG10(2) / "177464202324, "476747767461 / MCHD1680
C MCHD1700
C MACHINE CONSTANTS FOR PDP-11 FORTRAN'S SUPPORTING MCHD1720
C 32-BIT INTEGERS (EXPRESSED IN INTEGER AND OCTAL). MCHD1740
C MCHD1760
C DATA SMALL(1),SMALL(2) / 8388608, 0 / MCHD1780
C DATA LARGE(1),LARGE(2) / 2147483647, -1 / MCHD1800
C DATA RIGHT(1),RIGHT(2) / 612368384, 0 / MCHD1820
C DATA DIVER(1),DIVER(2) / 620756992, 0 / MCHD1840
C DATA LOG10(1),LOG10(2) / 1067065498, -2063872008 / MCHD1860
C MCHD1880
C DATA SMALL(1),SMALL(2) / 000040000000, 000000000000 / MCHD1900
C DATA LARGE(1),LARGE(2) / 017777777777, 037777777777 / MCHD1920
C DATA RIGHT(1),RIGHT(2) / 004440000000, 000000000000 / MCHD1940
C DATA DIVER(1),DIVER(2) / 004500000000, 000000000000 / MCHD1960
C DATA LOG10(1),LOG10(2) / 007746420232, 020476747770 / MCHD1980
C MCHD2000
C MACHINE CONSTANTS FOR PDP-11 FORTRAN'S SUPPORTING MCHD2020
C 16-BIT INTEGERS (EXPRESSED IN INTEGER AND OCTAL). MCHD2040
C MCHD2060
C DATA SMALL(1),SMALL(2) / 128, 0 / MCHD2080
C DATA SMALL(3),SMALL(4) / 0, 0 / MCHD2100
C MCHD2120
C DATA LARGE(1),LARGE(2) / 32767, -1 / MCHD2140
C DATA LARGE(3),LARGE(4) / -1, -1 / MCHD2160
C MCHD2180
C DATA RIGHT(1),RIGHT(2) / 9344, 0 / MCHD2200
C DATA RIGHT(3),RIGHT(4) / 0, 0 / MCHD2220
C MCHD2240
C DATA DIVER(1),DIVER(2) / 9472, 0 / MCHD2260
C DATA DIVER(3),DIVER(4) / 0, 0 / MCHD2280
C MCHD2300
C DATA LOG10(1),LOG10(2) / 16282, 8346 / MCHD2320
C DATA LOG10(3),LOG10(4) / -31493, -12296 / MCHD2340
C MCHD2360
C DATA SMALL(1),SMALL(2) / 0000200, 0000000 / MCHD2380
C DATA SMALL(3),SMALL(4) / 0000000, 0000000 / MCHD2400
C MCHD2420
C DATA LARGE(1),LARGE(2) / 0077777, 0177777 / MCHD2440
C DATA LARGE(3),LARGE(4) / 0177777, 0177777 / MCHD2460
C MCHD2480

```

```

C      DATA RIGHT(1),RIGHT(2) / 0022200, 0000000 / MCHD2500
C      DATA RIGHT(3),RIGHT(4) / 0000000, 0000000 / MCHD2520
C                                          MCHD2540
C      DATA DIVER(1),DIVER(2) / 0022400, 0000000 / MCHD2560
C      DATA DIVER(3),DIVER(4) / 0000000, 0000000 / MCHD2580
C                                          MCHD2600
C      DATA LOG10(1),LOG10(2) / 0037632, 0020232 / MCHD2620
C      DATA LOG10(3),LOG10(4) / 0102373, 0147770 / MCHD2640
C                                          MCHD2660
C      MACHINE CONSTANTS FOR THE UNIVAC 1100 SERIES. MCHD3700
C                                          MCHD3720
C      DATA SMALL(1),SMALL(2) / 0000040000000, 0000000000000 / MCHD3740
C      DATA LARGE(1),LARGE(2) / 0377777777777, 0777777777777 / MCHD3760
C      DATA RIGHT(1),RIGHT(2) / 0170540000000, 0000000000000 / MCHD3780
C      DATA DIVER(1),DIVER(2) / 0170640000000, 0000000000000 / MCHD3800
C      DATA LOG10(1),LOG10(2) / 0177746420232, 0411757177572 / MCHD3820
C                                          MCHD3860
C                                          MCHD4200
C      IF (I .LT. 1 .OR. I .GT. 5) MCHD4220
1      CALL SETERR(24HDMACH - I OUT OF BOUNDS,24,1,2) MCHD4240
C                                          MCHD4260
C      DIMACH = DMACH(I) MCHD4280
C      RETURN MCHD4300
C                                          MCHD4320
C      END MCHD4340

```

Automatic Error Handling

The second package provides a basic mechanism for dealing with the occurrence of errors.

In the PORT library [1], for which the package was developed, calls to the general subroutines in the library do not include flags for error indication in their calling sequences. Instead, when a called subroutine detects an error, it calls the principal error-handling routine, SETERR.

The package allows for two types of error, "fatal," and "recoverable," and a parameter in the call to SETERR must be set to specify the type. Fatal errors cause an error message to be printed, the run terminated, and a call made to a dump routine. (A dummy dump routine, FDUMP, is provided here.) For recoverable errors, unless the user has specifically requested to enter the recovery mode, similar events occur, an error message is printed, and the run terminated. Thus the process is failsafe for unwary users.

When the recovery mode is in effect, any call to SETERR given within a subprogram which has detected a recoverable error has the effect only of storing the fact that an error has occurred; the run is not terminated. The user, upon return from the subprogram, is responsible for testing for the occurrence of an error. If an error has occurred, the user must turn off the error state, because additional errors might arise and the occurrence of a recoverable error while in the error state constitutes an unrecoverable error, terminating the run.

Finally, since a called subprogram, say SUBA, may, in turn, call a lower-level subprogram containing recoverable errors, SUBA must check for the occurrence of errors in the lower-level routine and reinterpret them in the context of SUBA, which the user knows about. This means that SUBA must enter the recovery mode (saving the mode previously in effect), make the call to the lower-level subprogram, then, upon return from the lower-level routine, check for errors, and, before returning to the user, restore the previous recovery mode.

An error which has caused an invocation of SETERR has an associated number, message, and type (fatal or recoverable), and the effect of the error depends on whether the recovery mode is in effect or not. The various capabilities offered in the subprograms of the package are summarized as follows.

To signal that an error has occurred:

```
CALL SETERR(MESSG, NMESSG, NERR, IOPT)
```

where MESSG and NMESSG are, respectively, a Hollerith message and the number of characters in the message, and NERR is the error number. IOPT is used to specify the type of error: IOPT = 1 for a recoverable error, and IOPT = 2 for a fatal error.

To save the recovery (or nonrecovery) mode currently in effect, and enter a new one:

```
CALL ENTSRC(IROLD, IRNEW)
```

which saves the current mode in IROLD and sets the new one to IRNEW.

To avoid having multiple errors outstanding, it is a fatal error to call SETERR or ENTSRC if the error state is on, meaning that an error has occurred but has not been recovered from.

To restore the recovery (or nonrecovery) mode which was previously saved in IROLD:

```
CALL RETSRC(IROLD)
```

where RETSRC not only restores the previous mode, but also acts as a "safety" exit gate: Since multiple errors are illegal, RETSRC checks out the situation and allows return to the calling program only if (1) an error is not outstanding, or (2) the restored mode is recovery, so that the calling program is responsible for error checking.

To test if an error has occurred, and if its number was, say, 4, a statement such as the following is used:

```
IF (NERROR(NERR) .EQ. 4) GO TO 50
```

The value of the function NERROR and the value of the argument NERR are both set to the current value of the error number by NERROR. (The double assignment may be useful and comes free since Fortran prohibits functions with no arguments.) If the error number is nonzero, it means that an error has occurred and that corrective action must be taken.

To turn off the error state:

```
CALL ERROFF
```

In summary the user subprograms are:

SETERR	turns on the error state and saves a message and an error number;
ENTSRC	at entry, sets recovery (or nonrecovery) mode, provided no error state exists;
RETSRC	before returning, checks error situation and, if no errors exist, restores prior recovery (or nonrecovery) mode;
NERROR	returns the error number;
ERROFF	turns off the error state;
EPRINT	prints the error message.

These, in turn, call on the lower-level subprograms:

E9RINT	stores or prints error message, depending on switch setting;
S88FMT	sets up FORMAT array for printing;
I8SAVE	returns error number or recovery (or nonrecovery) mode, depending on one switch, and resets or does not reset the corresponding value depending on another;
FDUMP	a dummy routine to be replaced, if possible, by a locally written symbolic dump routine.

ALGORITHM

```

      SUBROUTINE SETERR(MESSG,NMESSG,NERR,IOPT)
C
C   SETERR SETS LERROR = NERR, OPTIONALLY PRINTS THE MESSAGE AND DUMPS
C   ACCORDING TO THE FOLLOWING RULES...
C
C   IF IOPT = 1 AND RECOVERING      - JUST REMEMBER THE ERROR.
C   IF IOPT = 1 AND NOT RECOVERING - PRINT AND STOP.
C   IF IOPT = 2                    - PRINT, DUMP AND STOP.
C
C INPUT
C
C MESSG - THE ERROR MESSAGE.
C NMESSG - THE LENGTH OF THE MESSAGE, IN CHARACTERS.
C NERR - THE ERROR NUMBER. MUST HAVE NERR NON-ZERO.
C IOPT - THE OPTION. MUST HAVE IOPT=1 OR 2.
C
C ERROR STATES -
C
C 1 - MESSAGE LENGTH NOT POSITIVE.
C 2 - CANNOT HAVE NERR=0.
C 3 - AN UNRECOVERED ERROR FOLLOWED BY ANOTHER ERROR.
C 4 - BAD VALUE FOR IOPT.
C
C ONLY THE FIRST 72 CHARACTERS OF THE MESSAGE ARE PRINTED.
C
C THE ERROR HANDLER CALLS A SUBROUTINE NAMED FDUMP TO PRODUCE A
C SYMBOLIC DUMP. TO COMPLETE THE PACKAGE, A DUMMY VERSION OF FDUMP
C IS SUPPLIED, BUT IT SHOULD BE REPLACED BY A LOCALLY WRITTEN VERSION
C WHICH AT LEAST GIVES A TRACE-BACK.
C
      INTEGER MESSG(1)
C
C THE UNIT FOR ERROR MESSAGES.
C
      IWUNIT=IIMACH(4)
C
      IF (NMESSG.GE.1) GO TO 10
C
C A MESSAGE OF NON-POSITIVE LENGTH IS FATAL.
C
      WRITE(IWUNIT,9000)
9000  FORMAT(52H1ERROR      1 IN SETERR - MESSAGE LENGTH NOT POSITIVE.)
      GO TO 60
C
C NW IS THE NUMBER OF WORDS THE MESSAGE OCCUPIES.
C
10  NW=(MIN0(NMESSG,72)-1)/IIMACH(6)+1
C
      IF (NERR.NE.0) GO TO 20
C
C CANNOT TURN THE ERROR STATE OFF USING SETERR.
C
      WRITE(IWUNIT,9001)
9001  FORMAT(42H1ERROR      2 IN SETERR - CANNOT HAVE NERR=0//
1      34H THE CURRENT ERROR MESSAGE FOLLOWS//)
      CALL EPRINT(MESSG,NW,NERR,.TRUE.)
      ITEMP=I8SAVE(1,1,.TRUE.)
      GO TO 50
C
C SET LERROR AND TEST FOR A PREVIOUS UNRECOVERED ERROR.
C
20  IF (I8SAVE(1,NERR,.TRUE.).EQ.0) GO TO 30
C
      WRITE(IWUNIT,9002)
9002  FORMAT(23H1ERROR      3 IN SETERR -,
1      48H AN UNRECOVERED ERROR FOLLOWED BY ANOTHER ERROR.//
2      48H THE PREVIOUS AND CURRENT ERROR MESSAGES FOLLOW.//)
      CALL EPRINT
      CALL EPRINT(MESSG,NW,NERR,.TRUE.)
      GO TO 50
C
C SAVE THIS MESSAGE IN CASE IT IS NOT RECOVERED FROM PROPERLY.

```

```

ERRS0000
ERRS0020
ERRS0040
ERRS0060
ERRS0080
ERRS0100
ERRS0120
ERRS0140
ERRS0160
ERRS0180
ERRS0200
ERRS0220
ERRS0240
ERRS0260
ERRS0280
ERRS0300
ERRS0320
ERRS0340
ERRS0360
ERRS0380
ERRS0400
ERRS0420
ERRS0440
ERRS0460
ERRS0480
ERRS0500
ERRS0520
ERRS0540
ERRS0560
ERRS0580
ERRS0600
ERRS0620
ERRS0640
ERRS0660
ERRS0680
ERRS0700
ERRS0720
ERRS0740
ERRS0760
ERRS0780
ERRS0800
ERRS0820
ERRS0840
ERRS0860
ERRS0880
ERRS0900
ERRS0920
ERRS0940
ERRS0960
ERRS0980
ERRS1000
ERRS1020
ERRS1040
ERRS1060
ERRS1080
ERRS1100
ERRS1120
ERRS1140
ERRS1160
ERRS1180
ERRS1200
ERRS1220
ERRS1240
ERRS1260
ERRS1280
ERRS1300
ERRS1320
ERRS1340
ERRS1360
ERRS1380
ERRS1400
ERRS1420

```



```

C
30 CALL E9RINT(MESSG,NW,NERR,.TRUE.)
C
      IF (IOPT.EQ.1 .OR. IOPT.EQ.2) GO TO 40
C
C MUST HAVE IOPT = 1 OR 2.
C
      WRITE(IWUNIT,9003)
9003  FORMAT(42H1ERROR      4 IN SETERR - BAD VALUE FOR IOPT//
1      34H THE CURRENT ERROR MESSAGE FOLLOWS//)
      GO TO 50
C
C TEST FOR RECOVERY.
C
40  IF (IOPT.EQ.2) GO TO 50
C
      IF (I8SAVE(2,0,.FALSE.).EQ.1) RETURN
C
      CALL EPRINT
      STOP
C
50  CALL EPRINT
60  CALL FDUMP
      STOP
C
      END

      ERRS1440
      ERRS1460
      ERRS1480
      ERRS1500
      ERRS1520
      ERRS1540
      ERRS1560
      ERRS1580
      ERRS1600
      ERRS1620
      ERRS1640
      ERRS1660
      ERRS1680
      ERRS1700
      ERRS1720
      ERRS1740
      ERRS1760
      ERRS1780
      ERRS1800
      ERRS1820
      ERRS1840
      ERRS1860
      ERRS1880
      ERRS1900
      ERRS1920
      ERRS1940

SUBROUTINE ENTSRC(IROLD,IRNEW)
C
C THIS ROUTINE RETURNS IROLD = LRECOV AND SETS LRECOV = IRNEW.
C
C IF THERE IS AN ACTIVE ERROR STATE, THE MESSAGE IS PRINTED
C AND EXECUTION STOPS.
C
C IRNEW = 0 LEAVES LRECOV UNCHANGED, WHILE
C IRNEW = 1 GIVES RECOVERY AND
C IRNEW = 2 TURNS RECOVERY OFF.
C
C ERROR STATES -
C
C 1 - ILLEGAL VALUE OF IRNEW.
C 2 - CALLED WHILE IN AN ERROR STATE.
C
      IF (IRNEW.LT.0 .OR. IRNEW.GT.2)
1      CALL SETERR(31HENTSRC - ILLEGAL VALUE OF IRNEW,31,1,2)
C
      IROLD=I8SAVE(2,IRNEW,IRNEW.NE.0)
C
C IF HAVE AN ERROR STATE, STOP EXECUTION.
C
      IF (I8SAVE(1,0,.FALSE.).NE.0) CALL SETERR
1      (39HENTSRC - CALLED WHILE IN AN ERROR STATE,39,2,2)
C
      RETURN
C
      END

      RECA0000
      RECA0020
      RECA0040
      RECA0060
      RECA0080
      RECA0100
      RECA0120
      RECA0140
      RECA0160
      RECA0180
      RECA0200
      RECA0220
      RECA0240
      RECA0260
      RECA0280
      RECA0300
      RECA0320
      RECA0340
      RECA0360
      RECA0380
      RECA0400
      RECA0420
      RECA0440
      RECA0460
      RECA0480
      RECA0500
      RECA0520
      RECA0540
      RECA0560

SUBROUTINE RETSRC(IROLD)
C
C THIS ROUTINE SETS LRECOV = IROLD.
C
C IF THE CURRENT ERROR BECOMES UNRECOVERABLE,
C THE MESSAGE IS PRINTED AND EXECUTION STOPS.
C
C ERROR STATES -
C
C 1 - ILLEGAL VALUE OF IROLD.
C
      IF (IROLD.LT.1 .OR. IROLD.GT.2)
1      CALL SETERR(31HRETSRC - ILLEGAL VALUE OF IROLD,31,1,2)
C
      ITEMP=I8SAVE(2,IROLD,.TRUE.)
C
C IF THE CURRENT ERROR IS NOW UNRECOVERABLE, PRINT AND STOP.
C
      RECB0000
      RECB0020
      RECB0040
      RECB0060
      RECB0080
      RECB0100
      RECB0120
      RECB0140
      RECB0160
      RECB0180
      RECB0200
      RECB0220
      RECB0240
      RECB0260
      RECB0280
      RECB0300
      RECB0320
      RECB0340

```

C	IF (IROLDEQ.1 .OR. I8SAVE(1,0,.FALSE.).EQ.0) RETURN	RECB0360
	CALL EPRINT	RECB0380
	STOP	RECB0400
C	END	RECB0420
		RECB0440
		RECB0460
	INTEGER FUNCTION NERR(NERR)	ERRN0000
C		ERRN0020
C	RETURNS NERR = NERR = THE VALUE OF THE ERROR FLAG LERROR.	ERRN0040
C		ERRN0060
	NERR=I8SAVE(1,0,.FALSE.)	ERRN0080
	NERR=NERR	ERRN0100
	RETURN	ERRN0120
C	END	ERRN0140
		ERRN0160
	SUBROUTINE ERROFF	ERRF0000
C		ERRF0020
C	URNS OFF THE ERROR STATE OFF BY SETTING LERROR=0.	ERRF0040
C		ERRF0060
	I=I8SAVE(1,0,.TRUE.)	ERRF0080
	RETURN	ERRF0100
C	END	ERRF0120
		ERRF0140
		ERRF0160
		ERRF0180
	SUBROUTINE EPRINT	ERRP0000
C		ERRP0020
C	THIS SUBROUTINE PRINTS THE LAST ERROR MESSAGE, IF ANY.	ERRP0040
C		ERRP0060
	INTEGER MESSG(1)	ERRP0080
C		ERRP0100
	CALL E9RINT(MESSG,1,1,.FALSE.)	ERRP0120
	RETURN	ERRP0140
C	END	ERRP0160
		ERRP0180
	SUBROUTINE E9RINT(MESSG,NW,NERR,SAVE)	ERRR0000
C		ERRR0020
C	THIS ROUTINE STORES THE CURRENT ERROR MESSAGE OR PRINTS THE OLD ONE,	ERRR0040
C	IF ANY, DEPENDING ON WHETHER OR NOT SAVE = .TRUE.	ERRR0060
C		ERRR0080
	INTEGER MESSG(NW)	ERRR0100
	LOGICAL SAVE	ERRR0120
C		ERRR0140
C	MESSGP STORES AT LEAST THE FIRST 72 CHARACTERS OF THE PREVIOUS	ERRR0160
C	MESSAGE. ITS LENGTH IS MACHINE DEPENDENT AND MUST BE AT LEAST	ERRR0180
C		ERRR0200
	1 + 71/(THE NUMBER OF CHARACTERS STORED PER INTEGER WORD).	ERRR0220
C		ERRR0240
	INTEGER MESSGP(36),FMT(14),CCPLUS	ERRR0260
C		ERRR0280
C	START WITH NO PREVIOUS MESSAGE.	ERRR0300
C		ERRR0320
	DATA MESSGP(1)/1H1/, NWP/0/, NERP/0/	ERRR0340
C		ERRR0360
C	SET UP THE FORMAT FOR PRINTING THE ERROR MESSAGE.	ERRR0380
C	THE FORMAT IS SIMPLY (A1,14X,72AXX) WHERE XX=IIMACH(6) IS THE	ERRR0400
C	NUMBER OF CHARACTERS STORED PER INTEGER WORD.	ERRR0420
C		ERRR0440
	DATA CCPLUS / 1H+ /	ERRR0460
C		ERRR0480
	DATA FMT(1) / 1H(/	ERRR0500
	DATA FMT(2) / 1HA /	ERRR0520
	DATA FMT(3) / 1H1 /	ERRR0540
	DATA FMT(4) / 1H, /	ERRR0560
	DATA FMT(5) / 1H1 /	ERRR0580
	DATA FMT(6) / 1H4 /	ERRR0600
	DATA FMT(7) / 1HX /	ERRR0620

```

DATA FMT( 8) / 1H, /
DATA FMT( 9) / 1H7 /
DATA FMT(10) / 1H2 /
DATA FMT(11) / 1HA /
DATA FMT(12) / 1HX /
DATA FMT(13) / 1HX /
DATA FMT(14) / 1H) /
C
IF (.NOT.SAVE) GO TO 20
C
C SAVE THE MESSAGE.
C
NWP=NW
NERRP=NERR
DO 10 I=1,NW
10 MESSGP(I)=MESSG(I)
C
GO TO 30
C
20 IF (I$SAVE(1,0,.FALSE.).EQ.0) GO TO 30
C
C PRINT THE MESSAGE.
C
IWUNIT=IIMACH(4)
WRITE(IWUNIT,9000) NERRP
9000 FORMAT(7H ERROR ,I4,4H IN )
C
CALL S88FMT(2,IIMACH(6),FMT(12))
WRITE(IWUNIT,FMT) CCPLUS,(MESSGP(I),I=1,NWP)
C
30 RETURN
C
END

SUBROUTINE S88FMT( N, W, IFMT )
C
C S88FMT REPLACES IFMT(1), ... , IFMT(N) WITH
C THE CHARACTERS CORRESPONDING TO THE N LEAST SIGNIFICANT
C DIGITS OF W.
C
INTEGER N,W,IFMT(N)
C
INTEGER NT,WT,DIGITS(10)
C
DATA DIGITS( 1) / 1H0 /
DATA DIGITS( 2) / 1H1 /
DATA DIGITS( 3) / 1H2 /
DATA DIGITS( 4) / 1H3 /
DATA DIGITS( 5) / 1H4 /
DATA DIGITS( 6) / 1H5 /
DATA DIGITS( 7) / 1H6 /
DATA DIGITS( 8) / 1H7 /
DATA DIGITS( 9) / 1H8 /
DATA DIGITS(10) / 1H9 /
C
NT = N
WT = W
C
10 IF (NT .LE. 0) RETURN
IDIGIT = MOD( WT, 10 )
IFMT(NT) = DIGITS(IDIGIT+1)
WT = WT/10
NT = NT - 1
GO TO 10
C
END

INTEGER FUNCTION I$SAVE(ISW,I$VALUE,SET)
C
C IF (ISW = 1) I$SAVE RETURNS THE CURRENT ERROR NUMBER AND
C SETS IT TO I$VALUE IF SET = .TRUE. .
C
C IF (ISW = 2) I$SAVE RETURNS THE CURRENT RECOVERY SWITCH AND
C SETS IT TO I$VALUE IF SET = .TRUE. .

```

```

ERRR0640
ERRR0660
ERRR0680
ERRR0700
ERRR0720
ERRR0740
ERRR0760
ERRR0780
ERRR0800
ERRR0820
ERRR0840
ERRR0860
ERRR0880
ERRR0900
ERRR0920
ERRR0940
ERRR0960
ERRR0980
ERRR1000
ERRR1020
ERRR1040
ERRR1060
ERRR1080
ERRR1100
ERRR1120
ERRR1140
ERRR1160
ERRR1180
ERRR1200
ERRR1220
ERRR1240
ERRR1260
ERRR1280

ERRM0000
ERRM0020
ERRM0040
ERRM0060
ERRM0080
ERRM0100
ERRM0120
ERRM0140
ERRM0160
ERRM0180
ERRM0200
ERRM0220
ERRM0240
ERRM0260
ERRM0280
ERRM0300
ERRM0320
ERRM0340
ERRM0360
ERRM0380
ERRM0400
ERRM0420
ERRM0440
ERRM0460
ERRM0480
ERRM0500
ERRM0520
ERRM0540
ERRM0560
ERRM0580
ERRM0600
ERRM0620

ERRV0000
ERRV0020
ERRV0040
ERRV0060
ERRV0080
ERRV0100
ERRV0120

```

```

C                                     ERRV0140
      LOGICAL SET                       ERRV0160
C                                     ERRV0180
      INTEGER IPARAM(2)                 ERRV0200
      EQUIVALENCE (IPARAM(1),LERROR) , (IPARAM(2),LRECOV) ERRV0220
C                                     ERRV0240
C START EXECUTION ERROR FREE AND WITH RECOVERY TURNED OFF. ERRV0260
C                                     ERRV0280
      DATA LERROR/0/ , LRECOV/2/      ERRV0300
C                                     ERRV0320
      I8SAVE=IPARAM(ISW)                ERRV0340
      IF (SET) IPARAM(ISW)=IVALUE       ERRV0360
C                                     ERRV0380
      RETURN                             ERRV0400
C                                     ERRV0420
      END                                 ERRV0440

      SUBROUTINE FDUMP                   FDMP0000
C THIS IS A DUMMY ROUTINE TO BE SENT OUT ON FDMP0020
C THE PORT SEDIT TAPE                  FDMP0040
C                                     FDMP0060
      RETURN                             FDMP0080
      END                                 FDMP0100

```

Dynamic Storage Allocation Using a Stack

The third package provides a basic mechanism for allocating and deallocating working storage on a storage stack.

In the PORT library [1], for which the package was developed, calls to the general subroutines in the library do not include, in their calling sequences, parameters representing scratch arrays; the work space is allocated and deallocated within the called subprograms.

To implement the stack in Fortran in a portable way, it has been put in labeled COMMON as a double-precision array of length 500:

```

COMMON/CSTAK/DSTAK(500)
DOUBLE PRECISION DSTAK

```

The stack handling capabilities, which are described more fully in [1], include stack allocation and deallocation (releasing space), stack initialization to a size different from the default length of 500 double-precision locations, stack query, i.e. the ability to find out dynamically how much space is available, modification of the length of the latest allocation, and finally, the ability to ascertain certain stack statistics, such as the number of outstanding allocations, the current active length, and the maximum active length achieved.

The various capabilities are summarized as follows:

To allocate (get) N locations of type ITYPE on the stack, set

```
INDEX = ISTKGT(N, ITYPE)
```

which returns an index into the stack for the first of the N items.

To deallocate (release) the last K allocations (not locations but entire allocations):

```
CALL ISTKRL(K)
```

To initialize the stack to, say, 1000 double-precision locations, use the subroutine, ISTKIN(N, ITYPE), as follows:

```

COMMON/CSTAK/DSTAK(1000)
DOUBLE PRECISION DSTAK
      ⋮
CALL ISTKIN(1000, 4)

```

To find out (query) how much space of type, ITYPE, is left:

NLEFT = ISTKQU(ITYPE)

To modify the length of the current outstanding allocation to N items:

INDEX = ISTKMD(N)

which will modify the length of the allocation to N items and, as in ISTKGT, return the index of the first item of that allocation.

To obtain certain stack statistics use the

INTEGER FUNCTION ISTKST(N)

In summary the user subprograms are:

ISTKGT allocates space on the stack;
 ISTKRL releases (deallocates) space;
 ISTKIN initializes (sets length) of stack;
 ISTKQU answers query as to space available on stack;
 ISTKMD modifies length of current allocation;
 ISTKST provides statistics on stack usage.

These, in turn, call on the lower-level subprogram:

I0TK00 initializes stack for special cases of nonstandard lengths for INTEGER, REAL, and DOUBLE PRECISION numbers.

ALGORITHM

```

      INTEGER FUNCTION ISTKGT(NITEMS, ITYPE)
C
C   ALLOCATES SPACE OUT OF THE INTEGER ARRAY ISTAK (IN COMMON
C   BLOCK CSTAK) FOR AN ARRAY OF LENGTH NITEMS AND OF TYPE
C   DETERMINED BY ITYPE AS FOLLOWS
C
C   1 - LOGICAL
C   2 - INTEGER
C   3 - REAL
C   4 - DOUBLE PRECISION
C   5 - COMPLEX
C
C   ON RETURN, THE ARRAY WILL OCCUPY
C
C   STAK(ISTKGT), STAK(ISTKGT+1), ..., STAK(ISTKGT-NITEMS+1)
C
C   WHERE STAK IS AN ARRAY OF TYPE ITYPE EQUIVALENCED TO ISTAK.
C
C   (FOR THOSE WANTING TO MAKE MACHINE DEPENDENT MODIFICATIONS
C   TO SUPPORT OTHER TYPES, CODES 6,7,8,9,10,11 AND 12 HAVE
C   BEEN RESERVED FOR 1/4 LOGICAL, 1/2 LOGICAL, 1/4 INTEGER,
C   1/2 INTEGER, QUAD PRECISION, DOUBLE COMPLEX AND QUAD
C   COMPLEX, RESPECTIVELY.)
C
C   THE ALLOCATOR RESERVES THE FIRST TEN INTEGER WORDS OF THE STACK
C   FOR ITS OWN INTERNAL BOOK-KEEPING. THESE ARE INITIALIZED BY
C   THE INITIALIZING SUBPROGRAM I0TK00 UPON THE FIRST CALL
C   TO A SUBPROGRAM IN THE ALLOCATION PACKAGE.
C
C   THE USE OF THE FIRST FIVE WORDS IS DESCRIBED BELOW.
C
C   ISTAK( 1) - LOUT, THE NUMBER OF CURRENT ALLOCATIONS.
C   ISTAK( 2) - LNOW, THE CURRENT ACTIVE LENGTH OF THE STACK.
C   ISTAK( 3) - LUSED, THE MAXIMUM VALUE OF ISTAK(2) ACHIEVED.
C   ISTAK( 4) - LMAX, THE MAXIMUM LENGTH THE STACK.
C   ISTAK( 5) - LBOOK, THE NUMBER OF WORDS USED FOR BOOKEEPING.
C
C   THE NEXT FIVE WORDS CONTAIN INTEGERS DESCRIBING THE AMOUNT
C   OF STORAGE ALLOCATED BY THE FORTRAN SYSTEM TO THE VARIOUS
C   DATA TYPES. THE UNIT OF MEASUREMENT IS ARBITRARY AND MAY
C   BE WORDS, BYTES OR BITS OR WHATEVER IS CONVENIENT. THE
      STKE0000
      STKE0020
      STKE0040
      STKE0060
      STKE0080
      STKE0100
      STKE0120
      STKE0140
      STKE0160
      STKE0180
      STKE0200
      STKE0220
      STKE0240
      STKE0260
      STKE0280
      STKE0300
      STKE0320
      STKE0340
      STKE0360
      STKE0380
      STKE0400
      STKE0420
      STKE0440
      STKE0460
      STKE0480
      STKE0500
      STKE0520
      STKE0540
      STKE0560
      STKE0580
      STKE0600
      STKE0620
      STKE0640
      STKE0660
      STKE0680
      STKE0700
      STKE0720
      STKE0740
      STKE0760
      STKE0780
      STKE0800

```



```

C ERROR STATES -
C
C 1 - NUMBER .LT. 0
C 2 - LNOW, LUSED, LMAX OR LBOOK OVERWRITTEN
C 3 - ATTEMPT TO DE-ALLOCATE NON-EXISTENT ALLOCATION
C 4 - THE POINTER AT ISTAK(LNOW) OVERWRITTEN
C
C COMMON /CSTAK/DSTAK
C
C DOUBLE PRECISION DSTAK(500)
C INTEGER ISTAK(1000)
C LOGICAL INIT
C
C EQUIVALENCE (DSTAK(1),ISTAK(1))
C EQUIVALENCE (ISTAK(1),LOUT)
C EQUIVALENCE (ISTAK(2),LNOW)
C EQUIVALENCE (ISTAK(3),LUSED)
C EQUIVALENCE (ISTAK(4),LMAX)
C EQUIVALENCE (ISTAK(5),LBOOK)
C
C DATA INIT/.TRUE./
C
C IF (INIT) CALL I0TK00(INIT,500,4)
C
C IF (NUMBER.LT.0) CALL SETERR(20HISTKRL - NUMBER.LT.0,20,1,2)
C
C IF (LNOW.LT.LBOOK.OR.LNOW.GT.LUSED.OR.LUSED.GT.LMAX) CALL SETERR
1 (47HISTKRL - LNOW, LUSED, LMAX OR LBOOK OVERWRITTEN,
2 47,2,2)
C
C IN = NUMBER
10 IF (IN.EQ.0) RETURN
C
C IF (LNOW.LE.LBOOK) CALL SETERR
1 (55HISTKRL - ATTEMPT TO DE-ALLOCATE NON-EXISTENT ALLOCATION,
2 55,3,2)
C
C CHECK TO MAKE SURE THE BACK POINTERS ARE MONOTONE.
C
C IF (ISTAK(LNOW).LT.LBOOK.OR.ISTAK(LNOW).GE.LNOW-1) CALL SETERR
1 (47HISTKRL - THE POINTER AT ISTAK(LNOW) OVERWRITTEN,
2 47,4,2)
C
C LOUT = LOUT-1
C LNOW = ISTAK(LNOW)
C IN = IN-1
C GO TO 10
C
C END

SUBROUTINE ISTKIN(NITEMS,ITYPE)
C
C INITIALIZES THE STACK ALLOCATOR, SETTING THE LENGTH OF THE STACK.
C
C ERROR STATES -
C
C 1 - NITEMS .LE. 0
C 2 - ITYPE .LE. 0 .OR. ITYPE .GE. 6
C
C LOGICAL INIT
C
C DATA INIT/.TRUE./
C
C IF (NITEMS.LE.0) CALL SETERR(20HISTKIN - NITEMS.LE.0,20,1,2)
C
C IF (ITYPE.LE.0.OR.ITYPE.GE.6) CALL SETERR
1 (33HISTKIN - ITYPE.LE.0.OR.ITYPE.GE.6,33,2,2)
C
C IF (INIT) CALL I0TK00(INIT,NITEMS,ITYPE)
C
C RETURN
C
C END

```

```

STKC0100
STKC0120
STKC0140
STKC0160
STKC0180
STKC0200
STKC0220
STKC0240
STKC0260
STKC0280
STKC0300
STKC0320
STKC0340
STKC0360
STKC0380
STKC0400
STKC0420
STKC0440
STKC0460
STKC0480
STKC0500
STKC0520
STKC0540
STKC0560
STKC0580
STKC0600
STKC0620
STKC0640
STKC0660
STKC0680
STKC0700
STKC0720
STKC0740
STKC0760
STKC0780
STKC0800
STKC0820
STKC0840
STKC0860
STKC0880
STKC0900
STKC0920
STKC0940
STKC0960
STKC0980
STKC1000
STKC1020
STKC1040
STKC1060

STKF0000
STKF0020
STKF0040
STKF0060
STKF0080
STKF0100
STKF0120
STKF0140
STKF0160
STKF0180
STKF0200
STKF0220
STKF0240
STKF0260
STKF0280
STKF0300
STKF0320
STKF0340
STKF0360
STKF0380
STKF0400
STKF0420
STKF0440

```

```

      INTEGER FUNCTION ISTKQU(ITYPE)
C
C RETURNS THE NUMBER OF ITEMS OF TYPE ITYPE THAT REMAIN
C TO BE ALLOCATED IN ONE REQUEST.
C
C ERROR STATES -
C
C 1 - LNOW, LUSED, LMAX OR LBOOK OVERWRITTEN
C 2 - ITYPE .LE. 0 .OR. ITYPE .GE. 6
C
C COMMON /CSTAK/DSTAK
C
C DOUBLE PRECISION DSTAK(500)
C INTEGER ISTAK(1000)
C INTEGER ISIZE(5)
C
C LOGICAL INIT
C
C EQUIVALENCE (DSTAK(1),ISTAK(1))
C EQUIVALENCE (ISTAK(2),LNOW)
C EQUIVALENCE (ISTAK(3),LUSED)
C EQUIVALENCE (ISTAK(4),LMAX)
C EQUIVALENCE (ISTAK(5),LBOOK)
C EQUIVALENCE (ISTAK(6),ISIZE(1))
C
C DATA INIT/.TRUE./
C
C IF (INIT) CALL I0TK00(INIT,500,4)
C
C IF (LNOW.LT.LBOOK.OR.LNOW.GT.LUSED.OR.LUSED.GT.LMAX) CALL SETERR
1 (47HISTKQU - LNOW, LUSED, LMAX OR LBOOK OVERWRITTEN,
2 47,1,2)
C
C IF (ITYPE.LE.0.OR.ITYPE.GE.6) CALL SETERR
1 (33HISTKQU - ITYPE.LE.0.OR.ITYPE.GE.6,33,2,2)
C
C ISTKQU = MAX0( ((LMAX-2)*ISIZE(2))/ISIZE(ITYPE)
1 - (LNOW*ISIZE(2)-1)/ISIZE(ITYPE)
2 - 1, 0 )
C
C RETURN
C
C END
      STKA0000
      STKA0020
      STKA0040
      STKA0060
      STKA0080
      STKA0100
      STKA0120
      STKA0140
      STKA0160
      STKA0180
      STKA0200
      STKA0220
      STKA0240
      STKA0260
      STKA0280
      STKA0300
      STKA0320
      STKA0340
      STKA0360
      STKA0380
      STKA0400
      STKA0420
      STKA0440
      STKA0460
      STKA0480
      STKA0500
      STKA0520
      STKA0540
      STKA0560
      STKA0580
      STKA0600
      STKA0620
      STKA0640
      STKA0660
      STKA0680
      STKA0700
      STKA0720
      STKA0740
      STKA0760
      STKA0780
      STKA0800
      STKA0820
      STKA0840

      INTEGER FUNCTION ISTKMD(NITEMS)
C
C CHANGES THE LENGTH OF THE FRAME AT THE TOP OF THE STACK
C TO NITEMS.
C
C ERROR STATES -
C
C 1 - LNOW OVERWRITTEN
C 2 - ISTAK(LNOWO-1) OVERWRITTEN
C
C COMMON /CSTAK/DSTAK
C
C DOUBLE PRECISION DSTAK(500)
C INTEGER ISTAK(1000)
C
C EQUIVALENCE (DSTAK(1),ISTAK(1))
C EQUIVALENCE (ISTAK(2),LNOW)
C
C LNOWO = LNOW
C CALL ISTKRL(1)
C
C ITYPE = ISTAK(LNOWO-1)
C
C IF (ITYPE.LE.0.OR.ITYPE.GE.6) CALL SETERR
1 (35HISTKMD - ISTAK(LNOWO-1) OVERWRITTEN,35,1,2)
C
C ISTKMD = ISTKGT(NITEMS,ITYPE)
C
C RETURN
C
C END
      STKB0000
      STKB0020
      STKB0040
      STKB0060
      STKB0080
      STKB0100
      STKB0120
      STKB0140
      STKB0160
      STKB0180
      STKB0200
      STKB0220
      STKB0240
      STKB0260
      STKB0280
      STKB0300
      STKB0320
      STKB0340
      STKB0360
      STKB0380
      STKB0400
      STKB0420
      STKB0440
      STKB0460
      STKB0480
      STKB0500
      STKB0520
      STKB0540
      STKB0560
      STKB0580
      STKB0600

```



```

      INTEGER FUNCTION IGTKST(NFACT)
C
C RETURNS CONTROL INFORMATION AS FOLLOWS
C
C NFACT      ITEM RETURNED
C
C   1          LOUT,  THE NUMBER OF CURRENT ALLOCATIONS
C   2          LNOW,  THE CURRENT ACTIVE LENGTH
C   3          LUSED, THE MAXIMUM USED
C   4          LMAX,  THE MAXIMUM ALLOWED
C
C      COMMON /CSTAK/DSTAK
C
C      DOUBLE PRECISION DSTAK(500)
C      INTEGER ISTAK(1000)
C      INTEGER ISTATS(4)
C      LOGICAL INIT
C
C      EQUIVALENCE (DSTAK(1),ISTAK(1))
C      EQUIVALENCE (ISTAK(1),ISTATS(1))
C
C      DATA INIT/.TRUE./
C
C      IF (INIT) CALL IGTK00(INIT,500,4)
C
C      IF (NFACT.LE.0.OR.NFACT.GE.5) CALL SETERR
C      1  (33HISTKST - NFACT.LE.0.OR.NFACT.GE.5,33,1,2)
C
C      ISTKST = ISTATS(NFACT)
C
C      RETURN
C
C      END
      STKS0000
      STKS0020
      STKS0040
      STKS0060
      STKS0080
      STKS0100
      STKS0120
      STKS0140
      STKS0160
      STKS0180
      STKS0200
      STKS0220
      STKS0240
      STKS0260
      STKS0280
      STKS0300
      STKS0320
      STKS0340
      STKS0360
      STKS0380
      STKS0400
      STKS0420
      STKS0440
      STKS0460
      STKS0480
      STKS0500
      STKS0520
      STKS0540
      STKS0560
      STKS0580
      STKS0600
      STKS0620
      STKS0640

      SUBROUTINE IGTK00(LARG,NITEMS,ITYPE)
C
C INITIALIZES THE STACK TO NITEMS OF TYPE ITYPE
C
C      COMMON /CSTAK/DSTAK
C
C      DOUBLE PRECISION DSTAK(500)
C      INTEGER ISTAK(1000)
C      LOGICAL LARG,INIT
C      INTEGER ISIZE(5)
C
C      EQUIVALENCE (DSTAK(1),ISTAK(1))
C      EQUIVALENCE (ISTAK(1),LOUT)
C      EQUIVALENCE (ISTAK(2),LNOW)
C      EQUIVALENCE (ISTAK(3),LUSED)
C      EQUIVALENCE (ISTAK(4),LMAX)
C      EQUIVALENCE (ISTAK(5),LBOOK)
C      EQUIVALENCE (ISTAK(6),ISIZE(1))
C
C      DATA INIT/.FALSE./
C
C      LARG = .FALSE.
C      IF (INIT) RETURN
C
C      HERE TO INITIALIZE
C
C      INIT = .TRUE.
C
C      SET DATA SIZES APPROPRIATE FOR A STANDARD CONFORMING
C      FORTRAN SYSTEM USING THE FORTRAN "STORAGE UNIT" AS THE
C      MEASURE OF SIZE.
C
C      LOGICAL
C      ISIZE(1) = 1
C      INTEGER
C      ISIZE(2) = 1
C      REAL
C      ISIZE(3) = 1
C      DOUBLE PRECISION
C      ISIZE(4) = 2
      STKG0000
      STKG0020
      STKG0040
      STKG0060
      STKG0080
      STKG0100
      STKG0120
      STKG0140
      STKG0160
      STKG0180
      STKG0200
      STKG0220
      STKG0240
      STKG0260
      STKG0280
      STKG0300
      STKG0320
      STKG0340
      STKG0360
      STKG0380
      STKG0400
      STKG0420
      STKG0440
      STKG0460
      STKG0480
      STKG0500
      STKG0520
      STKG0540
      STKG0560
      STKG0580
      STKG0600
      STKG0620
      STKG0640
      STKG0660
      STKG0680
      STKG0700
      STKG0720
      STKG0740
      STKG0760
      STKG0780

```

C	COMPLEX	STKG0800
	ISIZE(5) = 2	STKG0820
C		STKG0840
	LBOOK = 10	STKG0860
	LNOW = LBOOK	STKG0880
	LUSED = LBOOK	STKG0900
	LMAX = MAX0((NITEMS*ISIZE(ITYPE))/ISIZE(2), 12)	STKG0920
	LOUT = 0	STKG0940
C		STKG0960
	RETURN	STKG0980
C		STKG1000
	END	STKG1020

ACM Transactions on Mathematical Software, Vol. 5, No. 4, December 1979, Page 524.

REMARK ON ALGORITHM 528

Framework for a Portable Library [Z]

[P.A. Fox, A.D. Hall, and N.L. Schryer, *ACM Trans. Math. Software* 4, 2 (June 1978), 177-188]

Phyllis Fox [Recd 11 September 1979]

Bell Laboratories, 600 Mountain Ave., Murray Hill, NJ 07974

In the machine-dependent constants for the PDP 10 computer with KI processor, a transcription error has been found in the double-precision value for $\log_{10} b$.

The DATA line should read as follows:

C DATA LOG10(1),LOG10(2)/"177464202324,"047674776746/ MCHD1680

(See the complete listing of Algorithm 528, available from the ACM Algorithms Distribution Service or to be found in "Collected Algorithms from ACM.")

ACKNOWLEDGMENT

The author is grateful to Prof. Nelson Beebe of the Department of Physics, University of Utah, for detecting the incorrect value.

ALGORITHM 529

Permutations to Block Triangular Form [F1]

I. S. DUFF and J. K. REID
AERE Harwell

Key Words and Phrases: symmetric permutations, block triangular form, depth first search algorithm, sparse matrices
CR Categories: 5.0, 5.1, 5.3, 5.4
Language: ANSI Fortran

DESCRIPTION

Given the column numbers of the nonzeros in each row of a sparse matrix, this subroutine finds a symmetric permutation that makes the matrix block lower triangular. It can also be interpreted as accepting the row numbers of the nonzeros in each column and symmetrically permuting to block upper triangular form. If the user submits a matrix with zeros on the diagonal, subroutine MC13D might give a block triangular form which could be further reduced by unsymmetric permutations. To obtain the best results, the user is advised first to permute the matrix so that it has a zero-free diagonal. This can be done by Harwell subroutine MC21A (Duff [1]).

The subroutine is evoked by the Fortran statement

`CALL MC13D(N, ICN, LICN, IP, LENR, IOR, IB, NUM, IW)`

where the parameters are described in the listing given here.

The subroutine is based on Tarjan's depth first search algorithm [3], and its design is described in detail in Duff and Reid [2]. They also discuss experimental results from using this subroutine which has been tested on a wide range of both structured and randomly generated matrices.

This routine has been written in ANSI Fortran. Special comment cards have been included so that Harwell subroutine OE04A can be used to make an IBM Fortran version that uses half-length integers (INTEGER*2) for all the arrays except IP. This approximately halves the core requirements at the cost of restricting the order of the system to $2^{16} - 1$.

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2. DUFF, I.S., AND REID, J.K. An implementation of Tarjan's algorithm for the block triangularization of a matrix. *ACM Trans. Math. Software* 4, 2 (June 1978), 137-147.
3. TARJAN, R.E. Depth first search and linear graph algorithms. *SIAM J. Comput.* 1, pp. 146-160.

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ALGORITHM

[Subroutines MC13D and MC13E are printed here. A test deck for [MC13D is available from the ACM Distribution Service (see inside back cover for order form), or may be found in "Collected Algorithms from ACM."]

C I IS THE IBM VERSION	SI/MC1	10
C S IS THE STANDARD FORTRAN VERSION	MC1	20
C	MC1	30
SUBROUTINE MC13D(N, ICN, LICN, IP, LENR, IOR, IB, NUM, IW)	MC1	40
C	MC1	50
C DESCRIPTION OF PARAMETERS.	MC1	60
C INPUT VARIABLES N,ICN,LICN,IP,LENR.	MC1	70
C OUTPUT VARIABLES IOR,IB,NUM.	MC1	80
C	MC1	90
C N ORDER OF THE MATRIX.	MC1	100
C ICN ARRAY CONTAINING THE COLUMN INDICES OF THE NON-ZEROS. THOSE	MC1	110
C BELONGING TO A SINGLE ROW MUST BE CONTIGUOUS BUT THE ORDERING	MC1	120
C OF COLUMN INDICES WITHIN EACH ROW IS UNIMPORTANT AND WASTED	MC1	130
C SPACE BETWEEN ROWS IS PERMITTED.	MC1	140
C LICN LENGTH OF ARRAY ICN.	MC1	150
C IP IP(I), I=1,2,...N, IS THE POSITION IN ARRAY ICN OF THE FIRST	MC1	160
C COLUMN INDEX OF A NON-ZERO IN ROW I.	MC1	170
C LENR LENR(I) IS THE NUMBER OF NON-ZEROS IN ROW I, I=1,2,...N.	MC1	180
C IOR IOR(I) GIVES THE POSITION IN THE ORIGINAL ORDERING OF THE ROW	MC1	190
C OR COLUMN WHICH IS IN POSITION I IN THE PERMUTED FORM, I=1,2,...N.	MC1	200
C IB IB(I) IS THE ROW NUMBER IN THE PERMUTED MATRIX OF THE BEGINNING	MC1	210
C OF BLOCK I, I=1,2,...NUM.	MC1	220
C NUM NUMBER OF BLOCKS FOUND.	MC1	230
C IW WORK ARRAY .. SEE LATER COMMENTS.	MC1	240
C	MC1	250
INTEGER IP(N)	MC1	260
C INTEGER*2 ICN(LICN),LENR(N),IOR(N),IB(N),IW(N,3)	I/MC1	270
INTEGER ICN(LICN), LENR(N), IOR(N), IB(N), IW(N,3)	MC1	280
CALL MC13E(N, ICN, LICN, IP, LENR, IOR, IB, NUM, IW(1,1),	MC1	290
* IW(1,2), IW(1,3))	MC1	300
RETURN	MC1	310
END	MC1	320
SUBROUTINE MC13E(N, ICN, LICN, IP, LENR, ARP, IB, NUM, LOWL,	MC1	10
* NUMB, PREV)	MC1	20
INTEGER STP, DUMMY	MC1	30
INTEGER IP(N)	MC1	40
C	MC1	50
C ARP(I) IS ONE LESS THAN THE NUMBER OF UNSEARCHED EDGES LEAVING	MC1	60
C NODE I. AT THE END OF THE ALGORITHM IT IS SET TO A	MC1	70
C PERMUTATION WHICH PUTS THE MATRIX IN BLOCK LOWER	MC1	80
C TRIANGULAR FORM.	MC1	90
C IB(I) IS THE POSITION IN THE ORDERING OF THE START OF THE ITH	MC1	100
C BLOCK. IB(N+1-I) HOLDS THE NODE NUMBER OF THE ITH NODE	MC1	110
C ON THE STACK.	MC1	120
C LOWL(I) IS THE SMALLEST STACK POSITION OF ANY NODE TO WHICH A PATH	MC1	130
C FROM NODE I HAS BEEN FOUND. IT IS SET TO N+1 WHEN NODE I	MC1	140
C IS REMOVED FROM THE STACK.	MC1	150
C NUMB(I) IS THE POSITION OF NODE I IN THE STACK IF IT IS ON	MC1	160
C IT, IS THE PERMUTED ORDER OF NODE I FOR THOSE NODES	MC1	170
C WHOSE FINAL POSITION HAS BEEN FOUND AND IS OTHERWISE ZERO.	MC1	180
C PREV(I) IS THE NODE AT THE END OF THE PATH WHEN NODE I WAS	MC1	190
C PLACED ON THE STACK.	MC1	200
C INTEGER*2 ICN(LICN),LENR(N),ARP(N),IB(N),LOWL(N),NUMB(N),	I/MC1	210
C IPREV(N)	I/MC1	220
INTEGER ICN(LICN), LENR(N), ARP(N), IB(N), LOWL(N), NUMB(N),	MC1	230
* PREV(N)	MC1	240
C	MC1	250
C	MC1	260
C ICNT IS THE NUMBER OF NODES WHOSE POSITIONS IN FINAL ORDERING HAVE	MC1	270
C BEEN FOUND.	MC1	280
ICNT = 0	MC1	290
C NUM IS THE NUMBER OF BLOCKS THAT HAVE BEEN FOUND.	MC1	300
NUM = 0	MC1	310
NNM1 = N + N - 1	MC1	320
C	MC1	330
C INITIALIZATION OF ARRAYS.	MC1	340
DO 10 J=1,N	MC1	350
NUMB(J) = 0	MC1	360

ARP(J) = LENR(J) - 1	MC1 370
10 CONTINUE	MC1 380
C	MC1 390
C	MC1 400
DO 90 ISN=1,N	MC1 410
C LOOK FOR A STARTING NODE	MC1 420
IF (NUMB(ISN).NE.0) GO TO 90	MC1 430
IV = ISN	MC1 440
C IST IS THE NUMBER OF NODES ON THE STACK ... IT IS THE STACK POINTER.	MC1 450
IST = 1	MC1 460
C PUT NODE IV AT BEGINNING OF STACK.	MC1 470
LOWL(IV) = 1	MC1 480
NUMB(IV) = 1	MC1 490
IB(N) = IV	MC1 500
C	MC1 510
C THE BODY OF THIS LOOP PUTS A NEW NODE ON THE STACK OR BACKTRACKS.	MC1 520
DO 80 DUMMY=1,NNM1	MC1 530
I1 = ARP(IV)	MC1 540
C HAVE ALL EDGES LEAVING NODE IV BEEN SEARCHED.	MC1 550
IF (I1.LT.0) GO TO 30	MC1 560
I2 = IP(IV) + LENR(IV) - 1	MC1 570
I1 = I2 - I1	MC1 580
C	MC1 590
C LOOK AT EDGES LEAVING NODE IV UNTIL ONE ENTERS A NEW NODE OR	MC1 600
C ALL EDGES ARE EXHAUSTED.	MC1 610
DO 20 II=I1,I2	MC1 620
IW = ICN(II)	MC1 630
C HAS NODE IW BEEN ON STACK ALREADY.	MC1 640
IF (NUMB(IW).EQ.0) GO TO 70	MC1 650
C UPDATE VALUE OF LOWL(IV) IF NECESSARY.	MC1 660
IF (LOWL(IW).LT.LOWL(IV)) LOWL(IV) = LOWL(IW)	MC1 670
20 CONTINUE	MC1 680
C	MC1 690
C THERE ARE NO MORE EDGES LEAVING NODE IV.	MC1 700
ARP(IV) = -1	MC1 710
C IS NODE IV THE ROOT OF A BLOCK.	MC1 720
30 IF (LOWL(IV).LT.NUMB(IV)) GO TO 60	MC1 730
C	MC1 740
C ORDER NODES IN A BLOCK.	MC1 750
NUM = NUM + 1	MC1 760
IST1 = N + 1 - IST	MC1 770
LCNT = ICNT + 1	MC1 780
C PEEL BLOCK OFF THE TOP OF THE STACK STARTING AT THE TOP AND	MC1 790
C WORKING DOWN TO THE ROOT OF THE BLOCK.	MC1 800
DO 40 STP=IST1,N	MC1 810
IW = IB(STP)	MC1 820
LOWL(IW) = N + 1	MC1 830
ICNT = ICNT + 1	MC1 840
NUMB(IW) = ICNT	MC1 850
IF (IW.EQ.IV) GO TO 50	MC1 860
40 CONTINUE	MC1 870
50 IST = N - STP	MC1 880
IB(NUM) = LCNT	MC1 890
C ARE THERE ANY NODES LEFT ON THE STACK.	MC1 900
IF (IST.NE.0) GO TO 60	MC1 910
C HAVE ALL THE NODES BEEN ORDERED.	MC1 920
IF (ICNT.LT.N) GO TO 90	MC1 930
GO TO 100	MC1 940
C	MC1 950
C BACKTRACK TO PREVIOUS NODE ON PATH.	MC1 960
60 IW = IV	MC1 970
IV = PREV(IV)	MC1 980
C UPDATE VALUE OF LOWL(IV) IF NECESSARY.	MC1 990
IF (LOWL(IW).LT.LOWL(IV)) LOWL(IV) = LOWL(IW)	MC1 1000
GO TO 80	MC1 1010
C	MC1 1020
C PUT NEW NODE ON THE STACK.	MC1 1030
70 ARP(IV) = I2 - I1 - 1	MC1 1040
PREV(IW) = IV	MC1 1050
IV = IW	MC1 1060
IST = IST + 1	MC1 1070
LOWL(IV) = IST	MC1 1080
NUMB(IV) = IST	MC1 1090
K = N + 1 - IST	MC1 1100
IB(K) = IV	MC1 1110
80 CONTINUE	MC1 1120

C		MC1 1130
	90 CONTINUE	MC1 1140
C		MC1 1150
C		MC1 1160
C	PUT PERMUTATION IN THE REQUIRED FORM.	MC1 1170
	100 DO 110 I=1,N	MC1 1180
	II = NUMB(I)	MC1 1190
	ARP(II) = I	MC1 1200
	110 CONTINUE	MC1 1210
	RETURN	MC1 1220
	END	MC1 1230
C	TEST DECK FOR MC13D ... RUNS ON RANDOM MATRICES ...	MAN 10
	INTEGER IP(50), ICN(1000), IOR(50), IB(51), IW(150), LENR(50)	MAN 20
	INTEGER BLANK, EX, HOLD(100)	MAN 30
	LOGICAL A(50,50)	MAN 40
	DATA BLANK, EX, NOT /1H ,1HX,1H0/	MAN 50
	MM = 50	MAN 60
	LICN = 1000	MAN 70
C		MAN 80
C	MAIN LOOP.	MAN 90
	10 READ (5,99999) N, IPP	MAN 100
	IF (N.EQ.0) GO TO 100	MAN 110
	WRITE (6,99998) N, IPP	MAN 120
	DO 30 J=1,N	MAN 130
	DO 20 I=1,N	MAN 140
	A(I,J) = .FALSE.	MAN 150
	20 CONTINUE	MAN 160
	A(J,J) = .TRUE.	MAN 170
	30 CONTINUE	MAN 180
	IF (IPP.EQ.0) GO TO 60	MAN 190
	DO 50 K9=1,IPP	MAN 200
C	THESE STATEMENTS SHOULD BE REPLACED BY CALLS TO YOUR FAVOURITE	MAN 210
C	RANDOM NUMBER GENERATOR TO PLACE TWO PSEUDO-RANDOM NUMBERS	MAN 220
C	BETWEEN 1 AND N IN THE VARIABLES I AND J.	MAN 230
	40 CALL FA01BS(N, I)	MAN 240
	CALL FA01BS(N, J)	MAN 250
	IF (A(I,J)) GO TO 40	MAN 260
	A(I,J) = .TRUE.	MAN 270
	50 CONTINUE	MAN 280
C	SETUP CONVERTS MATRIX A(I,J) TO REQUIRED SPARSITY-ORIENTED	MAN 290
C	STORAGE FORMAT.	MAN 300
	60 CALL SETUP(N, A, MM, IP, ICN, LICN, LENR)	MAN 310
	CALL MC13D(N, ICN, LICN, IP, LENR, IOR, IB, NUM, IW)	MAN 320
C	OUTPUT REORDERED MATRIX WITH BLOCKING TO IMPROVE CLARITY.	MAN 330
	IF (NUM.EQ.1) WRITE (6,99997) NUM	MAN 340
	IF (NUM.NE.1) WRITE (6,99996) NUM	MAN 350
	IB(NUM+1) = N + 1	MAN 360
	INDEX = 100	MAN 370
	IBLOCK = 1	MAN 380
	DO 90 I=1,N	MAN 390
	DO 70 IJ=1,INDEX	MAN 400
	HOLD(IJ) = BLANK	MAN 410
	70 CONTINUE	MAN 420
	IF (I.EQ.IB(IBLOCK)) WRITE (6,99995)	MAN 430
	IF (I.EQ.IB(IBLOCK)) IBLOCK = IBLOCK + 1	MAN 440
	JBLOCK = 1	MAN 450
	INDEX = 0	MAN 460
	DO 80 J=1,N	MAN 470
	IF (J.EQ.IB(JBLOCK)) INDEX = INDEX + 1	MAN 480
	IF (J.EQ.IB(JBLOCK)) HOLD(INDEX) = BLANK	MAN 490
	IF (J.EQ.IB(JBLOCK)) JBLOCK = JBLOCK + 1	MAN 500
	INDEX = INDEX + 1	MAN 510
	II = IOR(I)	MAN 520
	JJ = IOR(J)	MAN 530
	IF (A(II,JJ)) HOLD(INDEX) = EX	MAN 540
	IF (.NOT.A(II,JJ)) HOLD(INDEX) = NOT	MAN 550
	80 CONTINUE	MAN 560
	WRITE (6,99994) HOLD	MAN 570
	90 CONTINUE	MAN 580
	WRITE (6,99993) (IB(I),I=1,NUM)	MAN 590
	GO TO 10	MAN 600
C		MAN 610
C	FORMAT STATEMENTS.	MAN 620
	100 STOP	MAN 630

99999	FORMAT (2I4)	MAN	640
99998	FORMAT (1H1, 20H MATRIX IS OF ORDER , I3, 9H AND HAS , I3, * 23H OFF-DIAGONAL NON-ZEROS)	MAN	650
99997	FORMAT (///31H THE REORDERED MATRIX WHICH HAS, I3, 10H BLOCK IS , * 11HOF THE FORM/)	MAN	660
99996	FORMAT (///31H THE REORDERED MATRIX WHICH HAS, I3, 10H BLOCKS IS, * 12H OF THE FORM/)	MAN	670
99995	FORMAT (3X)	MAN	680
99994	FORMAT (1X, 100A1)	MAN	690
99993	FORMAT (/46H THE STARTING POINT FOR EACH BLOCK IS GIVEN BY// * 20(2X, I4))	MAN	700
	END	MAN	710
		MAN	720
		MAN	730
		MAN	740
		MAN	750

	SUBROUTINE SETUP(N, A, MM, IP, ICN, LICN, LENR)	SET	10
	LOGICAL A(MM,MM)	SET	20
	INTEGER IP(N), ICN(LICN), LENR(N)	SET	30
	DO 10 I=1,N	SET	40
	LENR(I) = 0	SET	50
10	CONTINUE	SET	60
	IND = 1	SET	70
	DO 30 I=1,N	SET	80
	IP(I) = IND	SET	90
	DO 20 J=1,N	SET	100
	IF (.NOT.A(I,J)) GO TO 20	SET	110
	LENR(I) = LENR(I) + 1	SET	120
	ICN(IND) = J	SET	130
	IND = IND + 1	SET	140
20	CONTINUE	SET	150
30	CONTINUE	SET	160
	RETURN	SET	170
	END	SET	180

ALGORITHM 530

An Algorithm for Computing the Eigensystem of Skew-Symmetric Matrices and a Class of Symmetric Matrices [F2]

R. C. WARD and L. J. GRAY
Union Carbide Corporation, Nuclear Division

Key Words and Phrases: eigenvalues, eigenvectors, skew-symmetric matrices, symmetric matrices, matrices with constant diagonal

CR Categories: 5.14

Language: Fortran

DESCRIPTION

The set of Fortran subroutines given here is an implementation of the algorithm [3] for finding the eigenvectors x and eigenvalues λ such that $Ax = \lambda x$, where A is a real skew-symmetric matrix or a real tridiagonal symmetric matrix with a constant diagonal. The algorithm uses only orthogonal similarity transformations and is believed to be the most efficient procedure available for computing all the eigenvalues or the complete eigensystem for the indicated classes of matrices.

The three subroutines of the algorithm and their functions are described as follows:

TRIZD. A subroutine that transforms an arbitrary real skew-symmetric matrix to skew-symmetric tridiagonal form using orthogonal similarity transformations, saving the pertinent information about these transformations.

IMZD. A subroutine that computes the eigenvalues and, optionally, the eigenvectors of a symmetric tridiagonal matrix with zeros on the diagonal or of a skew-symmetric tridiagonal matrix.

TBAKZD. A subroutine that computes the eigenvectors of an arbitrary real skew-symmetric matrix by back-transforming the eigenvectors of the corresponding skew-symmetric tridiagonal matrix determined by TRIZD. Subroutines TRIZD and TBAKZD are straightforward adaptations of Fortran subroutines TRED1 and TRBAK1 [2] (originally published as Algol procedures in [1]) which accomplish similar functions for arbitrary real symmetric matrices. A detailed description of subroutine IMZD is given by Ward and Gray [3].

To determine the complete eigensystem of a full skew-symmetric matrix, the user should issue a call to TRIZD, IMZD, and TBAKZD, in that order. If only the eigenvalues are desired, then calling TBAKZD is unnecessary. To determine the complete eigensystem or only the eigenvalues of a tridiagonal skew-symmetric matrix or a tridiagonal symmetric matrix with zero diagonals, the user should only issue a call to IMZD. (If $\{\lambda_i\}$ and $\{x_i\}$ are the eigenvalues and eigenvectors of the matrix A , a tridiagonal symmetric matrix with zero diagonals, then the eigenvalues and eigenvectors of $A + \alpha I$ are $\{\lambda_i + \alpha\}$ and $\{x_i\}$.)

Subroutines TRIZD, IMZD, and TBAKZD have been tested extensively on

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the IBM 360/91 computer at Oak Ridge National Laboratory using double precision arithmetic. Some of the test cases are presented in [3]. Since only orthogonal similarity transformations are used, the algorithm is numerically stable, that is, each computed eigenvalue and its corresponding eigenvector are exact for a matrix close to the original matrix.

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ALGORITHM

```

C                                     MAIN0001
C ****                               MAIN0002
C                                     MAIN0003
C FUNCTION - THIS IS THE MAIN PROGRAM (DRIVER) FOR ILLUSTRATING THE MAIN0004
C           USE OF SUBROUTINES TRIZD, IMZD, AND TBAKZD. MAIN0005
C                                     MAIN0006
C REFERENCES - EIGENSYSTEM COMPUTATION FOR SKEW-SYMMETRIC MATRICES AND MAIN0007
C              A CLASS OF SYMMETRIC MATRICES, WARD,R C AND GRAY,L J. MAIN0008
C              TO APPEAR IN MANUSCRIPT SECTION OF ACM-TOMS. MAIN0009
C                                     MAIN0010
C              AN ALGORITHM FOR COMPUTING THE EIGENSYSTEM OF SKEW-
C              SYMMETRIC MATRICES AND A CLASS OF SYMMETRIC MATRICES, MAIN0012
C              TO APPEAR IN ALGORITHM SECTION OF ACM-TOMS. MAIN0013
C                                     MAIN0014
C REQUIRED FUNCTIONS FOR DRIVER AND SUBROUTINES - ABS,SIGN,SQRT,MOD MAIN0015
C                                     MAIN0016
C ****                               MAIN0017
C                                     MAIN0018
C      REAL  A(6,6),Z(6,6),E(6) MAIN0019
C      REAL  CON MAIN0020
C      INTEGER I,J,N,IM1,JP1,MAX,IERR MAIN0021
C      LOGICAL MATZ,SKEW MAIN0022
C      MAX = 6 MAIN0023
C      MATZ = .TRUE. MAIN0024
C                                     MAIN0025
C *** SET UP SKEW-SYMMETRIC TEST CASE MAIN0026
C                                     MAIN0027
C      WRITE (6,6000) MAIN0028
C 6000 FORMAT (1H1,47HEIGENSYSTEM COMPUTATION OF SKEW-SYMMETRIC TEST , MAIN0029
C           1 4HCASE//) MAIN0030
C      N = 5 MAIN0031
C      SKEW = .TRUE. MAIN0032
C                                     MAIN0033
C *** READ AND PRINT TEST MATRIX MAIN0034
C                                     MAIN0035
C      WRITE (6,6001) MAIN0036
C 6001 FORMAT (1H0,9X,11HTEST MATRIX/) MAIN0037
C      DO 10 I=1,N MAIN0038
C          READ (5,6002) (A(I,J),J=1,N) MAIN0039
C          WRITE (6,6002) (A(I,J),J=1,N) MAIN0040
C 6002 FORMAT (6F6.0) MAIN0041
C      10 CONTINUE MAIN0042
C                                     MAIN0043
C *** COMPUTE EIGENVALUES AND EIGENVECTORS MAIN0044
C                                     MAIN0045
C      CALL TRIZD (MAX,N,A,E) MAIN0046
C      CALL IMZD (N,E,MATZ,SKEW,MAX,Z,IERR) MAIN0047
C      IF (IERR .NE. 0) WRITE (6,6003) IERR MAIN0048
C 6003 FORMAT (1H0/1H0,11HIMZD IERR =,15) MAIN0049
C      CALL TBAKZD (MAX,N,A,N,MAX,Z) MAIN0050
C                                     MAIN0051
C *** PRINT EIGENVALUES AND EIGENVECTORS MAIN0052
C                                     MAIN0053
C      WRITE (6,6004) MAIN0054
C 6004 FORMAT (1H0/1H0,2X,11HEIGENVALUES,25X,12HEIGENVECTORS) MAIN0055
C      J = 0 MAIN0056
C      15 J = J + 1 MAIN0057

```

```

        WRITE (6,6005)                                MAIN0058
6005    FORMAT (/)                                    MAIN0059
        IF (E(J) .EQ. 0.D0) GO TO 20                    MAIN0060
        JP1 = J+1                                       MAIN0061
        WRITE (6,6006) E(J),Z(1,J),Z(1,JP1)           MAIN0062
6006    FORMAT (1X,E15.8,4H * I,5X,E15.8,5H + ,E15.8,4H * I) MAIN0063
        WRITE (6,6007) (Z(I,J),Z(I,JP1),I=2,N)       MAIN0064
6007    FORMAT (25X,E15.8,5H + ,E15.8,4H * I)       MAIN0065
        WRITE (6,6005)                                MAIN0066
        CON = -Z(1,JP1)                                MAIN0067
        WRITE (6,6006) E(JP1),Z(1,J),CON             MAIN0068
        DO 18 I=2,N                                    MAIN0069
            CON = -Z(I,JP1)                            MAIN0070
            WRITE (6,6007) Z(I,J),CON                 MAIN0071
18    CONTINUE                                       MAIN0072
        J = J + 1                                       MAIN0073
        GO TO 30                                       MAIN0074
20    WRITE (6,6008) E(J),Z(1,J)                     MAIN0075
6008    FORMAT (1X,E15.8,9X,E15.8)                 MAIN0076
        WRITE (6,6009) (Z(I,J),I=2,N)               MAIN0077
6009    FORMAT (25X,E15.8)                          MAIN0078
        30 IF (J .LT. N) GO TO 15                    MAIN0079
C                                                    MAIN0080
C *** SET UP TRIDIAGONAL, SYMMETRIC, ZERO DIAGONAL TEST CASE MAIN0081
C                                                    MAIN0082
        WRITE (6,6010)                                MAIN0083
6010    FORMAT (1H1,50HEIGENSYSTEM COMPUTATION OF TRIDIAGONAL, SYMMETRIC, ,MAIN0084
1      24H ZERO DIAGONAL TEST CASE//)              MAIN0085
        N = 6                                          MAIN0086
        SKEW = .FALSE.                                MAIN0087
        DO 40 I=1,N                                    MAIN0088
            E(I) = 1.                                  MAIN0089
            DO 40 J=1,N                                MAIN0090
                A(I,J) = 0.                            MAIN0091
40    CONTINUE                                       MAIN0092
        DO 50 I=2,N                                    MAIN0093
            IM1 = I-1                                  MAIN0094
            A(I,IM1) = E(I)                            MAIN0095
            A(IM1,I) = E(I)                            MAIN0096
50    CONTINUE                                       MAIN0097
C                                                    MAIN0098
C *** PRINT TEST MATRIX                             MAIN0099
C                                                    MAIN0100
        WRITE (6,6011)                                MAIN0101
6011    FORMAT (1H0,12X,11HTEST MATRIX/)           MAIN0102
        DO 60 I=1,N                                    MAIN0103
            WRITE (6,6002) (A(I,J),J=1,N)             MAIN0104
60    CONTINUE                                       MAIN0105
C                                                    MAIN0106
C *** COMPUTE EIGENVALUES AND EIGENVECTORS          MAIN0107
C                                                    MAIN0108
        CALL IMZD (N,E,MATZ,SKEW,MAX,Z,IERR)         MAIN0109
        IF (IERR .NE. 0) WRITE (6,6003) IERR        MAIN0110
C                                                    MAIN0111
C *** PRINT EIGENVALUES AND EIGENVECTORS           MAIN0112
C                                                    MAIN0113
        WRITE (6,6012)                                MAIN0114
6012    FORMAT (1H0/1H0,2X,11HEIGENVALUES,13X,12HEIGENVECTORS) MAIN0115
        DO 70 J=1,N                                    MAIN0116
            WRITE (6,6005)                             MAIN0117
            WRITE (6,6008) E(J),Z(1,J)                 MAIN0118
            WRITE (6,6009) (Z(I,J),I=2,N)             MAIN0119
70    CONTINUE                                       MAIN0120
        STOP                                          MAIN0121
        END                                          MAIN0122

        SUBROUTINE TRIZD ( NA, N, A, E )              TRIZ0001
C                                                    TRIZ0002
C ****                                             TRIZ0003
C                                                    TRIZ0004
C FUNCTION - REDUCES A REAL SKEW-SYMMETRIC MATRIX TO A SKEW-SYMMETRIC TRIZ0005
C            TRIDIAGONAL MATRIX USING ORTHOGONAL SIMILARITY TRIZ0006
C            TRANSFORMATIONS                          TRIZ0007
C                                                    TRIZ0008
C PARAMETERS                                         TRIZ0009

```

```

C
C   NA      - INPUT INTEGER SPECIFYING THE ROW DIMENSION OF A AS      TRIZ0011
C             DECLARED IN THE CALLING PROGRAM DIMENSION STATEMENT    TRIZ0012
C
C   N       - INPUT INTEGER SPECIFYING THE ORDER OF A                 TRIZ0013
C
C   A(NA,N) - ON INPUT, A CONTAINS THE REAL SKEW-SYMMETRIC MATRIX.   TRIZ0014
C             ONLY THE STRICT LOWER TRIANGLE OF THE MATRIX NEED      TRIZ0015
C             BE SUPPLIED.                                           TRIZ0016
C             ON OUTPUT, A CONTAINS INFORMATION ABOUT THE ORTHOGONAL  TRIZ0017
C             TRANSFORMATIONS USED IN THE REDUCTION IN ITS FULL      TRIZ0018
C             LOWER TRIANGLE. THE STRICT UPPER TRIANGLE OF A IS      TRIZ0019
C             UNALTERED.                                             TRIZ0020
C
C   E(N)    - OUTPUT ARRAY CONTAINING THE LOWER SUBDIAGONAL ELEMENTS TRIZ0021
C             OF THE TRIDIAGONAL MATRIX IN ITS LAST N-1 POSITIONS.  TRIZ0022
C             E(1) IS SET TO ZERO.                                    TRIZ0023
C
C   REQUIRED FUNCTIONS - ABS,SIGN,SQRT                                TRIZ0024
C
C   ****                                         TRIZ0025
C
C   REAL    A(NA,N),E(N)                                         TRIZ0026
C   REAL    F,G,H,SCALE                                         TRIZ0027
C   REAL    ABS,SIGN,SQRT                                       TRIZ0028
C   INTEGER I,J,K,L,IL,JM1,JP1                                  TRIZ0029
C
C   IF (N .EQ. 1) GO TO 230                                       TRIZ0030
C
C   *** MAIN DO LOOP I=N STEP -1 UNTIL 2                            TRIZ0031
C
C   DO 220 II = 2, N                                               TRIZ0032
C     I = N + 2 - II                                               TRIZ0033
C     L = I - 1                                                    TRIZ0034
C     H = 0.                                                       TRIZ0035
C     SCALE = 0.                                                  TRIZ0036
C
C   *** NORMALIZE ROW                                             TRIZ0037
C
C     DO 100 K = 1, L                                             TRIZ0038
C       SCALE = SCALE + ABS(A(I,K))                               TRIZ0039
C     100 CONTINUE                                               TRIZ0040
C
C     IF (SCALE .NE. 0.) GO TO 120                                TRIZ0041
C     E(I) = 0.                                                  TRIZ0042
C     GO TO 215                                                  TRIZ0043
C
C   *** COMPUTE ELEMENTS OF U VECTOR                               TRIZ0044
C
C     120 DO 130 K = 1, L                                         TRIZ0045
C         A(I,K) = A(I,K) / SCALE                                TRIZ0046
C         H = H + A(I,K) * A(I,K)                               TRIZ0047
C     130 CONTINUE                                               TRIZ0048
C
C     F = A(I,L)                                                 TRIZ0049
C     G = -SIGN(SQRT(H),F)                                       TRIZ0050
C     E(I) = SCALE * G                                          TRIZ0051
C     H = H - F * G                                             TRIZ0052
C     A(I,L) = F - G                                            TRIZ0053
C     IF (L .EQ. 1) GO TO 200                                    TRIZ0054
C
C   *** COMPUTE ELEMENTS OF A*U/H                                 TRIZ0055
C
C     DO 180 J = 1, L                                             TRIZ0056
C       G = 0.                                                  TRIZ0057
C       IF (J .EQ. 1) GO TO 150                                  TRIZ0058
C       JM1 = J - 1                                             TRIZ0059
C
C     DO 140 K = 1, JM1                                          TRIZ0060
C       G = G + A(J,K) * A(I,K)                                  TRIZ0061
C     140 CONTINUE                                               TRIZ0062
C
C     150 IF (J .EQ. L) GO TO 170                                TRIZ0063
C         JP1 = J + 1                                           TRIZ0064
C
C     DO 160 K = JP1, L                                         TRIZ0065

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          G = G - A(K,J) * A(I,K)
160      CONTINUE
C
170      E(J) = G / H
180      CONTINUE
C
C *** COMPUTE REDUCED A
C
      DO 190 J = 2, L
          F = A(I,J)
          G = E(J)
          JM1 = J - 1
C
          DO 190 K = 1, JM1
              A(J,K) = A(J,K) + F * E(K) - G * A(I,K)
190      CONTINUE
C
200      DO 210 K = 1, L
          A(I,K) = SCALE * A(I,K)
210      CONTINUE
C
215      A(I,I) = SCALE * SQRT(H)
220      CONTINUE
C
230      E(1) = 0.
          RETURN
          END
          TRIZ0086
          TRIZ0087
          TRIZ0088
          TRIZ0089
          TRIZ0090
          TRIZ0091
          TRIZ0092
          TRIZ0093
          TRIZ0094
          TRIZ0095
          TRIZ0096
          TRIZ0097
          TRIZ0098
          TRIZ0099
          TRIZ0100
          TRIZ0101
          TRIZ0102
          TRIZ0103
          TRIZ0104
          TRIZ0105
          TRIZ0106
          TRIZ0107
          TRIZ0108
          TRIZ0109
          TRIZ0110
          TRIZ0111
          TRIZ0112

          SUBROUTINE IMZD ( N, E, MATZ, SKEW, NZ, Z, IERR )
C
C ****
C
C FUNCTION - COMPUTE THE EIGENVALUES AND OPTIONALLY THE EIGENVECTORS
C           OF A SYMMETRIC TRIDIAGONAL MATRIX WITH ZERO DIAGONALS
C           OR A SKEW-SYMMETRIC TRIDIAGONAL MATRIX USING AN
C           IMPLICIT QR-TYPE ITERATION
C
C PARAMETERS
C
C   N      - INPUT INTEGER SPECIFYING THE ORDER OF THE TRIDIAGONAL
C           MATRIX
C
C   E(N)   - ON INPUT, ARRAY CONTAINING THE LOWER SUBDIAGONAL
C           ELEMENTS OF THE TRIDIAGONAL MATRIX IN ITS LAST N-1
C           POSITIONS. E(1) IS ARBITRARY.
C           ON OUTPUT, ARRAY CONTAINS THE EIGENVALUES. THE NON-ZERO
C           EIGENVALUES OCCUR IN PAIRS WITH OPPOSITE SIGNS AND
C           ARE FOUND IN ADJACENT LOCATIONS IN E. THE EIGENVALUES
C           OF SYMMETRIC MATRICES ARE REAL AND THE EIGENVALUES
C           OF SKEW-SYMMETRIC MATRICES ARE PURELY IMAGINARY
C           COMPLEX NUMBERS. IF AN ERROR EXIT IS MADE, THE
C           EIGENVALUES ARE CORRECT FOR INDICES IERR+1,IERR+2...
C
C   MATZ   - INPUT LOGICAL VARIABLE SPECIFYING THE EIGENVECTOR
C           OPTION
C           = .TRUE.  EIGENVECTORS ARE TO BE COMPUTED
C           = .FALSE. EIGENVECTORS ARE NOT TO BE COMPUTED
C
C   SKEW   - INPUT LOGICAL VARIABLE SPECIFYING TYPE OF INPUT MATRIX
C           = .TRUE.  INPUT TRIDIAGONAL MATRIX IS SKEW-SYMMETRIC
C           = .FALSE. INPUT TRIDIAGONAL MATRIX IS SYMMETRIC WITH
C           ZERO DIAGONALS
C           SKEW IS NOT REFERENCED IF MATZ = .FALSE.
C
C   NZ     - INPUT INTEGER SPECIFYING THE ROW DIMENSION OF Z AS
C           DECLARED IN THE CALLING PROGRAM DIMENSION STATEMENT
C
C   Z(NZ,N) - OUTPUT ARRAY CONTAINING THE ORTHOGONAL EIGENVECTORS
C           OF THE INPUT TRIDIAGONAL MATRIX. EIGENVECTORS CORRE-
C           SPONDING TO ZERO EIGENVALUES ARE NORMALIZE TO UNIT
C           2-NORM (LENGTH) AND THOSE CORRESPONDING TO NON-ZERO
C           EIGENVALUES HAVE 2-NORM OF SQUARE ROOT 2. IF THE J-TH
C           EIGENVALUE IS ZERO OR REAL (I.E. E(J)), ITS EIGEN-
C           VECTOR IS FOUND IN THE J-TH COLUMN OF Z. IF THE J-TH
C           EIGENVALUE IS IMAGINARY (I.E. E(J)*I) WITH E(J+1) =
          IMZD0001
          IMZD0002
          IMZD0003
          IMZD0004
          IMZD0005
          IMZD0006
          IMZD0007
          IMZD0008
          IMZD0009
          IMZD0010
          IMZD0011
          IMZD0012
          IMZD0013
          IMZD0014
          IMZD0015
          IMZD0016
          IMZD0017
          IMZD0018
          IMZD0019
          IMZD0020
          IMZD0021
          IMZD0022
          IMZD0023
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          IMZD0028
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          IMZD0030
          IMZD0031
          IMZD0032
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          IMZD0038
          IMZD0039
          IMZD0040
          IMZD0041
          IMZD0042
          IMZD0043
          IMZD0044
          IMZD0045
          IMZD0046
          IMZD0047

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C          -E(J), THE REAL PART OF ITS EIGENVECTOR IS FOUND IN IMZD0048
C          THE J-TH COLUMN OF Z AND ITS IMAGINARY PART FOUND IN IMZD0049
C          THE (J+1)-TH COLUMN. IF AN ERROR EXIT IS MADE, Z IMZD0050
C          CONTAINS THE EIGENVECTORS ASSOCIATED WITH THE STORED IMZD0051
C          EIGENVALUES. IMZD0052
C          Z IS NOT REFERENCED IF MATZ = .FALSE. IMZD0053
C
C          IEERR - OUTPUT ERROR CODE IMZD0054
C          = 0 NORMAL RETURN (ALL EIGENVALUES/VECTORS FOUND) IMZD0056
C          = J IF THE J-TH EIGENVALUE HAS NOT BEEN DETERMINED IMZD0057
C              AFTER 30 ITERATIONS IMZD0058
C          IMZD0059
C  REQUIRED FUNCTIONS - ABS,SIGN,SQRT,MOD IMZD0060
C
C ***** IMZD0061
C
C          REAL E(N),Z(NZ,N) IMZD0064
C          REAL F,G,Q,C,S,R,P,TEST,TMAG IMZD0065
C          REAL ABS,SIGN,SQRT IMZD0066
C          INTEGER I,J,K,L,M,LO,LOM1,MO,LS,IM1,JP1,KM1,KP1,LM1,MM1, IMZD0067
C          I IP3,IEERR,IEO,ITS,IP1 IMZD0068
C          INTEGER MOD IMZD0069
C          LOGICAL MATZ,SKEW IMZD0070
C
C          IF (.NOT. MATZ) GO TO 115 IMZD0071
C
C          *** PLACE IDENTITY MATRIX IN Z IMZD0072
C
C          DO 110 I = 1, N IMZD0073
C              DO 100 J = 1, N IMZD0074
C                  Z(I,J) = 0. IMZD0075
C          100 CONTINUE IMZD0076
C              Z(I,I) = 1. IMZD0077
C          110 CONTINUE IMZD0078
C
C          115 IEERR = 0 IMZD0079
C              M = N IMZD0080
C              MM1 = M - 1 IMZD0081
C              E(1) = 0. IMZD0082
C              ITS = 0 IMZD0083
C
C          120 IF (M .LT. 2) GO TO 370 IMZD0084
C              MO = M IMZD0085
C
C          *** SEARCH FOR NEXT SUBMATRIX TO SOLVE (MATRIX SPLITTING) IMZD0086
C
C          F = 0. IMZD0087
C          DO 130 I = 1, MM1 IMZD0088
C              J = M - I IMZD0089
C              JP1 = J + 1 IMZD0090
C              G = ABS(E(JP1)) IMZD0091
C              TMAG = ABS(E(J)) + F IMZD0092
C              TEST = TMAG + G IMZD0093
C              IF (TEST .EQ. TMAG) GO TO 140 IMZD0094
C              F = G IMZD0095
C          130 CONTINUE IMZD0096
C              JP1 = 1 IMZD0097
C
C          140 LO = JP1 + 1 IMZD0098
C              LOM1 = JP1 IMZD0099
C              IF (LOM1 .EQ. M) GO TO 290 IMZD0100
C              IF (.NOT. MATZ) GO TO 160 IMZD0101
C              IF (.NOT. SKEW) GO TO 160 IMZD0102
C
C          C ***** PLACE CORRECT SIGN ON IDENTITY DIAGONALS IMZD0103
C
C          DO 150 I = LOM1, M, 4 IMZD0104
C              Z(I,I) = -Z(I,I) IMZD0105
C              IP3 = I + 3 IMZD0106
C              IF (IP3 .GT. M) GO TO 160 IMZD0107
C              Z(IP3,IP3) = -Z(IP3,IP3) IMZD0108
C          150 CONTINUE IMZD0109
C
C          160 IF (LO .EQ. M) GO TO 300 IMZD0110
C              IEO = M - LO IMZD0111
C              IEO = MOD(IEO,2) IMZD0112

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```

      L = L0
      IF (IEO .EQ. 0) GO TO 230
C
C *** FIND ZERO EIGENVALUE OF ODD ORDERED SUBMATRICES
C
      C = 0.
      S = -1.
      DO 190 I = L0, MM1, 2
        K = MM1 + L0 - I
        KP1 = K + 1
        Q = -S * E(KP1)
        E(KP1) = C * E(KP1)
        IF (ABS(E(K)) .GT. ABS(Q)) GO TO 170
        C = E(K) / Q
        R = SQRT(C*C + 1.)
        E(K) = Q * R
        S = 1. / R
        C = C * S
        GO TO 180
      170 S = Q / E(K)
        R = SQRT(1. + S*S)
        E(K) = E(K) * R
        C = 1. / R
        S = S * C
      180 IF (.NOT. MATZ) GO TO 190
C
C *** ACCUMULATE TRANSFORMATIONS FOR EIGENVECTORS
C
      KM1 = K - 1
      Z(KM1,M) = -S * Z(KM1,KM1)
      Z(KM1,KM1) = C * Z(KM1,KM1)
      DO 185 J = KP1, M, 2
        Z(J,KM1) = S * Z(J,M)
        Z(J,M) = C * Z(J,M)
      185 CONTINUE
C
      190 CONTINUE
      M = MM1
      MM1 = M - 1
      IF (L0 .EQ. M) GO TO 300
C
C *** CHECK FOR CONVERGENCE OR SMALL SUBDIAGONAL ELEMENT
C
      200 DO 210 I = L0, MM1, 2
        K = MM1 + L0 - I
        L = K + 1
        TMAG = ABS(E(L)) + ABS(E(K-1))
        TEST = TMAG + E(K)
        IF (TEST .EQ. TMAG) GO TO 220
      210 CONTINUE
      L = L0
      220 IF (L .EQ. M) GO TO 300
C
C *** FORM SHIFT
C
      230 ITS = ITS + 1
      IF (ITS .GT. 30) GO TO 360
      F = E(M-3)
      G = E(M-2)
      C = E(MM1)
      S = E(M)
      P = ((C-F) * (C+F) + (S-G) * (S+G)) / (2. * G * C)
      R = SQRT(P*P + 1.)
      Q = (G / (P + SIGN(R,P))) - C
      F = E(L)
      LM1 = L - 1
      E(LM1) = ((F-S) * (F+S) + C * Q) / F
C
C *** PERFORM ONE IMPLICIT QR ITERATION ON CHOLESKY FACTOR
C
      LS = LOM1
      C = 1.
      S = 1.
      DO 280 I = L, MM1
        IP1 = I + 1
        IM1 = I - 1

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IMZD0124
 IMZD0125
 IMZD0126
 IMZD0127
 IMZD0128
 IMZD0129
 IMZD0130
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 IMZD0198
 IMZD0199

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      Q = S * E(IP1)
      E(IP1) = C * E(IP1)
      IF (ABS(E(IM1)) .GT. ABS(Q)) GO TO 240
      C = E(IM1) / Q
      R = SQRT(C*C + 1.)
      E(IM1) = Q * R
      S = 1. / R
      C = C * S
      GO TO 250
240   S = Q / E(IM1)
      R = SQRT(1. + S*S)
      E(IM1) = E(IM1) * R
      C = 1. / R
      S = S * C
250   F = E(IP1)
      E(IP1) = -S * E(I) + C * F
      E(I) = C * E(I) + S * F
      IF (.NOT. MATZ) GO TO 280
C
C *** ACCUMULATE TRANSFORMATIONS FOR EIGENVECTORS
C
      DO 260 J = LS, MO, 2
        F = Z(J,IP1)
        Z(J,IP1) = -S * Z(J,IM1) + C * F
        Z(J,IM1) = C * Z(J,IM1) + S * F
260   CONTINUE
      IF (LS .EQ. LOM1) GO TO 270
      LS = LOM1
      GO TO 280
270   LS = L0
280   CONTINUE
      E(LM1) = 0.
      GO TO 200
C
C *** ITERATION CONVERGED TO ONE ZERO EIGENVALUE
C
290   E(M) = 0.
      M = MM1
      GO TO 310
C
C *** ITERATION CONVERGED TO EIGENVALUE PAIR
C
300   E(MM1) = E(M)
      E(M) = -E(M)
      M = M - 2
C
310   ITS = 0
      MM1 = M - 1
      IF (M .GT. L0) GO TO 200
      IF (M .EQ. L0) GO TO 300
      IF (.NOT. MATZ) GO TO 120
      IF (SKEW) GO TO 120
C
C *** COMPUTE EIGENVECTORS FROM ORTHONORMAL COLUMNS OF Z IF NOT SKEW
C
320   K = M0
330   IF (E(K) .EQ. 0.) GO TO 350
      KM1 = K - 1
      DO 340 J = LOM1, M0, 2
        Z(J,K) = Z(J,KM1)
        F = Z(J+1,K)
        Z(J+1,KM1) = F
        Z(J+1,K) = -F
340   CONTINUE
      K = KM1
350   K = K-1
      IF (K .GT. LOM1) GO TO 330
      IF (IERR .NE. 0) GO TO 370
      GO TO 120
C
C *** ERROR EXIT
C
360   IERR = M
      IF (.NOT. MATZ) GO TO 370
      IF (.NOT. SKEW) GO TO 320
C
370   RETURN
      END

```

IMZD0200
 IMZD0201
 IMZD0202
 IMZD0203
 IMZD0204
 IMZD0205
 IMZD0206
 IMZD0207
 IMZD0208
 IMZD0209
 IMZD0210
 IMZD0211
 IMZD0212
 IMZD0213
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 IMZD0274
 IMZD0275
 IMZD0276
 IMZD0277


```

      SUBROUTINE TBAKZD ( NA, N, A, M, NZ, Z )
C
C *****
C
C FUNCTION - FORMS THE EIGENVECTORS OF A REAL SKEW-SYMMETRIC MATRIX
C             BY BACK TRANSFORMING THOSE OF THE CORRESPONDING SKEW-
C             SYMMETRIC TRIDIAGONAL MATRIX DETERMINED BY TRIZD
C
C PARAMETERS
C
C   NA      - INPUT INTEGER SPECIFYING THE ROW DIMENSION OF A
C             AS DECLARED IN CALLING PROGRAM DIMENSION STATEMENT
C
C   N       - INPUT INTEGER SPECIFYING THE ORDER OF A
C
C   A(NA,N) - INPUT ARRAY CONTAINING INFORMATION ABOUT THE ORTHOGONAL
C             TRANSFORMATIONS USED IN THE REDUCTION BY TRIZD IN
C             ITS FULL LOWER TRIANGLE
C
C   M       - INPUT INTEGER SPECIFYING THE NUMBER OF EIGENVECTORS TO
C             BE BACK TRANSFORMED
C
C   NZ      - INPUT INTEGER SPECIFYING THE ROW DIMENSION OF Z AS
C             DECLARED IN CALLING PROGRAM DIMENSION STATEMENT
C
C   Z(NZ,M) - ON INPUT, Z CONTAINS THE REAL AND IMAGINARY (IF
C             COMPLEX) PARTS OF THE EIGENVECTORS TO BE BACK TRANS-
C             FORMED IN ITS FIRST M COLUMNS
C             ON OUTPUT, Z CONTAINS THE REAL AND IMAGINARY (IF
C             COMPLEX) PARTS OF THE TRANSFORMED EIGENVECTORS IN
C             ITS FIRST M COLUMNS
C *****
C
C   REAL  A(NA,N),Z(NZ,M)
C   REAL  H,S
C   INTEGER I,J,K,L
C
C   IF (M .EQ. 0) GO TO 140
C   IF (N .EQ. 1) GO TO 140
C
C   DO 130 I = 2, N
C     L = I - 1
C     H = A(I,I)
C     IF (H .EQ. 0.) GO TO 130
C
C     DO 120 J = 1, M
C       S = 0.
C
C       DO 100 K = 1, L
C         S = S + A(I,K) * Z(K,J)
C100    CONTINUE
C
C       S = (S / H) / H
C
C       DO 110 K = 1, L
C         Z(K,J) = Z(K,J) - S * A(I,K)
C110    CONTINUE
C
C120    CONTINUE
C
C130    CONTINUE
C
C140    RETURN
C      END
C
C      0.   2.  -2.   0.  -4.
C     -2.   0.   4.  -5.   3.
C      2.  -4.   0.   1.   1.
C      0.   5.  -1.   0.   2.
C      4.  -3.  -1.  -2.   0.

```

TBAK0001
TBAK0002
TBAK0003
TBAK0004
TBAK0005
TBAK0006
TBAK0007
TBAK0008
TBAK0009
TBAK0010
TBAK0011
TBAK0012
TBAK0013
TBAK0014
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TBAK0055
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TBAK0057
TBAK0058
TBAK0059
TBAK0060
TBAK0061
TBAK0062
TBAK0063
TBAK0064
TBAK0065

ALGORITHM 531

Contour Plotting [J6]

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Key Words and Phrases: contour plotting
CR Categories: 8.2
Language: Fortran

DESCRIPTION

Given a two-dimensional array of samples of a surface, contour values, and a subroutine to draw lines, the subroutine GCONTR determines sequences of points in the plane which may be used to draw contours through equal values of the surface.

A contour plotting algorithm may be constructed by following contours from some starting point until they either close or intersect a boundary, or by examining each cell of the grid in turn and drawing all contours found inside the cell. The advantages of contour following are that contour labeling is relatively easy and that the pen of an incremental plotter does not move about as much without writing anything. The advantages of methods which draw all contours found inside a cell are that less auxiliary storage is needed and that each cell can be completely processed before going on to the next, thereby allowing generation of contours over a much larger array than can be accommodated in main memory at one time.

GCONTR is of the type which follows contours. On a representative problem, a program using GCONTR generated about 11,000 plotter commands with about 57,000 commands generated by a program using a cellular method. For this problem, grid lines, user identification, and a table of contour values required about 36,000 additional commands.

Since the data used by contour plotting programs are presented at discrete points and since the contour values are not necessarily equal to the dependent variable values at the nodal points of the mesh, some method must be used to estimate the point at which a contour value intersects the edge of a cell. If one uses nonlinear interpolation, there is the possibility of multiple intersections of a contour with an edge and the possibility of closed contours which intersect no edges; if one wants to follow the interpolated surface faithfully, drawing the contour entails finding zeros of a nonlinear bivariate function. Linear interpolation

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This work was sponsored by the National Aeronautics and Space Administration under Contract NAS 7-100; this paper presents the results of one phase of research carried out at the Jet Propulsion Laboratory. The work was carried out using a Univac 1108 computer, CalComp model 565 drum plotters, a Calcomp model 1675 Com, and Tektronix 4000 series interactive graphics terminals.

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has the drawback that extrema occur only at nodal points. GCONTR uses linear interpolation. If this does not produce sufficiently smooth, accurate contours, we recommend computing more function values or interpolatory subtabulation to a finer mesh. If the required storage is larger than desired, the independent variable plane is easily divided into separately processed segments. GCONTR does not provide automatic subtabulation because different methods work better for different problems. Since GCONTR uses a contour following method, any implementation of automatic subtabulation would necessarily have a higher cost than a reasonable user implementation.

Once the points of intersection of a contour with edges have been determined, one must still decide how to draw the contour line between such intersections. We recommend that straight lines be used. Most other curves can cause balloons and crossing of contours of different values.

We now discuss the features which distinguish GCONTR from the methods of Cottafava and Le Moli [1] and Crane [2]. GCONTR and the algorithms described in [1] and [2] are similar in that all are contour following methods, the data are assumed to be presented on the nodes of a topologically rectangular grid, and linear interpolation is used along the edges of a cell.

GCONTR and algorithm [2] begin following a contour as soon as an edge of a cell of the coordinate grid is found which the contour line intersects. Edges intersected by the contour are flagged in an integer array in [2] and by marks in a bitmap in GCONTR. In algorithm [1] a preliminary detection of all intersections of edges and contours is done and then the contours are drawn. The details of the bookkeeping are not described in [1].

Algorithms [1] and [2] start the contour at the first edge found while searching along rows or columns of the data array. GCONTR reduces pen-up travel by searching along a rectangular spiral starting at the current pen position. All three algorithms draw all contours which intersect boundaries before drawing any contours which do not.

Algorithm [1] treats the case of a contour intersecting all four edges of a cell in a way which depends on the fixed order of examining the edges. If a cell is approached from different directions, different contours will be drawn. Algorithm [2] forces the contours to cross inside the cell. When several different contours intersect all four edges, the resulting pattern looks like an asterisk. GCONTR examines the linear interpolate along the "top" and "bottom" edges of the cell. If the interpolate on the top edge is less than the interpolate on the bottom edge, the contour line is drawn top-to-left and bottom-to-right. Otherwise, the contour line is drawn top-to-right and bottom-to-left. This method selects the same contours if the axes are exchanged. To show this, note that the use of linear interpolation along the edges of the unit cell implies that there is some estimated function value F_x which occurs at the same abscissa on the top and bottom edges, and some estimated function value F_y which occurs at the same ordinate on the "left" and "right" edges. It is easy to show that $F_x = F_y$. Therefore, the contour of value F_x is the only contour which crosses inside the cell, and it divides the cell into four rectangular regions. In two of these regions the estimated function value is less than F_x . In the other two, the estimated function value is greater than F_x .

Algorithms [1] and [2] consider all data in the array. GCONTR provides a means whereby the user may designate excluded elements of the array. The edges connecting excluded elements to elements not excluded are not examined. This feature has been found to be very useful in practice, as data are frequently not presented on a complete rectangular grid. When using methods which do not have this feature, one frequently finds extraneous contours drawn between the regions containing data and the regions containing no data.

The output method is not described in [1]. The output from algorithm [2] is an ordered set of points in an array which define the contour lines to be drawn. GCONTR calls a line drawing subroutine provided by the user. If the user wishes to take the risks noted above associated with using other than straight lines for


```

NWORD = (N-1)/NBPW           IGE 130
NBIT = MOD(N-1,NBPW)         IGE 140
IGET = MOD(BITMAP(NWORD+1)/2**(NBPW-NBIT-1),2) IGE 150
RETURN                        IGE 160
END                            IGE 170

SUBROUTINE GCONTR(Z, NRZ, NX, NY, CV, MCV, ZMAX, BITMAP, DRAW) GCO 10
C                               GCO 20
C THIS SUBROUTINE DRAWS A CONTOUR THROUGH EQUAL VALUES OF AN ARRAY. GCO 30
C                               GCO 40
C ***** FORMAL ARGUMENTS ***** GCO 50
C                               GCO 60
C Z IS THE ARRAY FOR WHICH CONTOURS ARE TO BE DRAWN. THE ELEMENTS GCO 70
C OF Z ARE ASSUMED TO LIE UPON THE NODES OF A TOPOLOGICALLY GCO 80
C RECTANGULAR COORDINATE SYSTEM - E.G. CARTESIAN, POLAR (EXCEPT GCO 90
C THE ORIGIN), ETC. GCO 100
C                               GCO 110
C NRZ IS THE NUMBER OF ROWS DECLARED FOR Z IN THE CALLING PROGRAM. GCO 120
C                               GCO 130
C NX IS THE LIMIT FOR THE FIRST SUBSCRIPT OF Z. GCO 140
C                               GCO 150
C NY IS THE LIMIT FOR THE SECOND SUBSCRIPT OF Z. GCO 160
C                               GCO 170
C CV ARE THE VALUES OF THE CONTOURS TO BE DRAWN. GCO 180
C                               GCO 190
C MCV IS THE NUMBER OF CONTOUR VALUES IN CV. GCO 200
C                               GCO 210
C ZMAX IS THE MAXIMUM VALUE OF Z FOR CONSIDERATION. A VALUE OF GCO 220
C Z(I,J) GREATER THAN ZMAX IS A SIGNAL THAT THAT POINT AND THE GCO 230
C GRID LINE SEGMENTS RADIATING FROM THAT POINT TO IT'S NEIGHBORS GCO 240
C ARE TO BE EXCLUDED FROM CONTOURING. GCO 250
C                               GCO 260
C BITMAP IS A WORK AREA LARGE ENOUGH TO HOLD 2*NX*NY*NCV BITS. IT GCO 270
C IS ACCESSED BY LOW-LEVEL ROUTINES, WHICH ARE DESCRIBED BELOW. GCO 280
C LET J BE THE NUMBER OF USEFUL BITS IN EACH WORD OF BITMAP, GCO 290
C AS DETERMINED BY THE USER MACHINE AND IMPLEMENTATION OF GCO 300
C THE BITMAP MANIPULATION SUBPROGRAMS DESCRIBED BELOW. THEN GCO 310
C THE NUMBER OF WORDS REQUIRED FOR THE BITMAP IS THE FLOOR OF GCO 320
C (2*NX*NY*NCV+J-1)/J. GCO 330
C                               GCO 340
C DRAW IS A USER-PROVIDED SUBROUTINE USED TO DRAW CONTOURS. GCO 350
C THE CALLING SEQUENCE FOR DRAW IS: GCO 360
C                               GCO 370
C CALL DRAW (X,Y,IFLAG) GCO 380
C LET NX = INTEGER PART OF X, FX = FRACTIONAL PART OF X. GCO 390
C THEN X SHOULD BE INTERPRETED SUCH THAT INCREASES IN NX GCO 400
C CORRESPOND TO INCREASES IN THE FIRST SUBSCRIPT OF Z, AND GCO 410
C FX IS THE FRACTIONAL DISTANCE FROM THE ABCISSA CORRESPONDING GCO 420
C TO NX TO THE ABCISSA CORRESPONDING TO NX+1, GCO 430
C AND Y SHOULD BE INTERPRETED SIMILARLY FOR THE SECOND GCO 440
C SUBSCRIPT OF Z. GCO 450
C THE LOW-ORDER DIGIT OF IFLAG WILL HAVE ONE OF THE VALUES: GCO 460
C 1 - CONTINUE A CONTOUR, GCO 470
C 2 - START A CONTOUR AT A BOUNDARY, GCO 480
C 3 - START A CONTOUR NOT AT A BOUNDARY, GCO 490
C 4 - FINISH A CONTOUR AT A BOUNDARY, GCO 500
C 5 - FINISH A CLOSED CONTOUR (NOT AT A BOUNDARY). GCO 510
C NOTE THAT REQUESTS 1, 4 AND 5 ARE FOR PEN-DOWN GCO 520
C MOVES, AND THAT REQUESTS 2 AND 3 ARE FOR PEN-UP GCO 530
C MOVES. GCO 540
C 6 - SET X AND Y TO THE APPROXIMATE 'PEN' POSITION, USING GCO 550
C THE NOTATION DISCUSSED ABOVE. THIS CALL MAY BE GCO 560
C IGNORED, THE RESULT BEING THAT THE 'PEN' POSITION GCO 570
C IS TAKEN TO CORRESPOND TO Z(1,1). GCO 580
C IFLAG/10 IS THE CONTOUR NUMBER. GCO 590
C                               GCO 600
C ***** EXTERNAL SUBPROGRAMS ***** GCO 610
C                               GCO 620
C DRAW IS THE USER-SUPPLIED LINE DRAWING SUBPROGRAM DESCRIBED ABOVE. GCO 630
C DRAW MAY BE SENSITIVE TO THE HOST COMPUTER AND TO THE PLOT DEVICE. GCO 640
C FILLO IS USED TO FILL A BITMAP WITH ZEROES. CALL FILLO (BITMAP,N) GCO 650
C FILLS THE FIRST N BITS OF BITMAP WITH ZEROES. GCO 660
C MARK1 IS USED TO PLACE A 1 IN A SPECIFIC BIT OF THE BITMAP. GCO 670
C CALL MARK1 (BITMAP,N) PUTS A 1 IN THE NTH BIT OF THE BITMAP. GCO 680

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C      IGET IS USED TO DETERMINE THE SETTING OF A PARTICULAR BIT IN THE GCO 690
C      BITMAP. I=IGET(BITMAP,N) SETS I TO ZERO IF THE NTH BIT OF THE GCO 700
C      BITMAP IS ZERO, AND SETS I TO ONE IF THE NTH BIT IS ONE. GCO 710
C      FILLO, MARK1 AND IGET ARE MACHINE SENSITIVE. GCO 720
C GCO 730
C *****GCO 740
C GCO 750
C      REAL Z(NRZ,1), CV(1) GCO 760
C      INTEGER BITMAP(1) GCO 770
C      INTEGER L1(4), L2(4), IJ(2) GCO 780
C GCO 790
C      L1 AND L2 CONTAIN LIMITS USED DURING THE SPIRAL SEARCH FOR THE GCO 800
C      BEGINNING OF A CONTOUR. GCO 810
C      IJ STORES SUBSCRIPTS USED DURING THE SPIRAL SEARCH. GCO 820
C GCO 830
C      INTEGER I1(2), I2(2), I3(6) GCO 840
C GCO 850
C      I1, I2 AND I3 ARE USED FOR SUBSCRIPT COMPUTATIONS DURING THE GCO 860
C      EXAMINATION OF LINES FROM Z(I,J) TO IT'S NEIGHBORS. GCO 870
C GCO 880
C      REAL XINT(4) GCO 890
C GCO 900
C      XINT IS USED TO MARK INTERSECTIONS OF THE CONTOUR UNDER GCO 910
C      CONSIDERATION WITH THE EDGES OF THE CELL BEING EXAMINED. GCO 920
C GCO 930
C      REAL XY(2) GCO 940
C GCO 950
C      XY IS USED TO COMPUTE COORDINATES FOR THE DRAW SUBROUTINE.. GCO 960
C GCO 970
C      EQUIVALENCE (L2(1),IMAX), (L2(2),JMAX), (L2(3),IMIN), GCO 980
C      * (L2(4),JMIN) GCO 990
C      EQUIVALENCE (IJ(1),I), (IJ(2),J) GCO 1000
C      EQUIVALENCE (XY(1),X), (XY(2),Y) GCO 1010
C GCO 1020
C      DATA L1(3) /-1/, L1(4) /-1/ GCO 1030
C      DATA I1 /1,0/, I2 /1,-1/, I3 /1,0,0,1,1,0/ GCO 1040
C GCO 1050
C      L1(1) = NX GCO 1060
C      L1(2) = NY GCO 1070
C      DMAX = ZMAX GCO 1080
C GCO 1090
C      SET THE CURRENT PEN POSITION. THE DEFAULT POSITION CORRESPONDS GCO 1100
C      TO Z(1,1). GCO 1110
C GCO 1120
C      X = 1.0 GCO 1130
C      Y = 1.0 GCO 1140
C      CALL DRAW(X, Y, 6) GCO 1150
C      ICUR = MAXO(1,MINO(INT(X),NX)) GCO 1160
C      JCUR = MAXO(1,MINO(INT(Y),NY)) GCO 1170
C GCO 1180
C      CLEAR THE BITMAP GCO 1190
C GCO 1200
C      CALL FILLO(BITMAP, 2*NX*NY*NCV) GCO 1210
C GCO 1220
C      SEARCH ALONG A RECTANGULAR SPIRAL PATH FOR A LINE SEGMENT HAVING GCO 1230
C      THE FOLLOWING PROPERTIES: GCO 1240
C      1. THE END POINTS ARE NOT EXCLUDED, GCO 1250
C      2. NO MARK HAS BEEN RECORDED FOR THE SEGMENT, GCO 1260
C      3. THE VALUES OF Z AT THE ENDS OF THE SEGMENT ARE SUCH THAT GCO 1270
C      ONE Z IS LESS THAN THE CURRENT CONTOUR VALUE, AND THE GCO 1280
C      OTHER IS GREATER THAN OR EQUAL TO THE CURRENT CONTOUR GCO 1290
C      VALUE. GCO 1300
C GCO 1310
C      SEARCH ALL BOUNDARIES FIRST, THEN SEARCH INTERIOR LINE SEGMENTS. GCO 1320
C      NOTE THAT THE INTERIOR LINE SEGMENTS NEAR EXCLUDED POINTS MAY BE GCO 1330
C      BOUNDARIES. GCO 1340
C GCO 1350
C      IBKEY = 0 GCO 1360
C      10 I = ICUR GCO 1370
C      J = JCUR GCO 1380
C      20 IMAX = I GCO 1390
C      IMIN = -I GCO 1400
C      JMAX = J GCO 1410
C      JMIN = -J GCO 1420
C      IDIR = 0 GCO 1430
C      DIRECTION ZERO IS +I, 1 IS +J, 2 IS -I, 3 IS -J. GCO 1440

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```

30 NXIDIR = IDIR + 1
K = NXIDIR
IF (NXIDIR.GT.3) NXIDIR = 0
40 I = IABS(I)
J = IABS(J)
IF (Z(I,J).GT.DMAX) GO TO 140
L = 1
C L=1 MEANS HORIZONTAL LINE, L=2 MEANS VERTICAL LINE.
50 IF (IJ(L).GE.L1(L)) GO TO 130
II = I + I1(L)
JJ = J + I1(3-L)
IF (Z(II,JJ).GT.DMAX) GO TO 130
ASSIGN 100 TO JUMP
C THE NEXT 15 STATEMENTS (OR SO) DETECT BOUNDARIES.
60 IX = 1
IF (IJ(3-L).EQ.1) GO TO 80
II = I - I1(3-L)
JJ = J - I1(L)
IF (Z(II,JJ).GT.DMAX) GO TO 70
II = I + I2(L)
JJ = J + I2(3-L)
IF (Z(II,JJ).LT.DMAX) IX = 0
70 IF (IJ(3-L).GE.L1(3-L)) GO TO 90
80 II = I + I1(3-L)
JJ = J + I1(L)
IF (Z(II,JJ).GT.DMAX) GO TO 90
IF (Z(I+1,J+1).LT.DMAX) GO TO JUMP, (100, 280)
90 IX = IX + 2
GO TO JUMP, (100, 280)
100 IF (IX.EQ.3) GO TO 130
IF (IX+IBKEY.EQ.0) GO TO 130
C NOW DETERMINE WHETHER THE LINE SEGMENT IS CROSSED BY THE CONTOUR.
II = I + I1(L)
JJ = J + I1(3-L)
Z1 = Z(I,J)
Z2 = Z(II,JJ)
DO 120 ICV=1,NCV
IF (IGET(BITMAP,2*(NX*(NY*(ICV-1)+J-1)+I-1)+L).NE.0) GO TO 120
IF (CV(ICV).LE.AMIN1(Z1,Z2)) GO TO 110
IF (CV(ICV).LE.AMAX1(Z1,Z2)) GO TO 190
110 CALL MARK1(BITMAP, 2*(NX*(NY*(ICV-1)+J-1)+I-1)+L)
120 CONTINUE
130 L = L + 1
IF (L.LE.2) GO TO 50
140 L = MOD(IDIR,2) + 1
IJ(L) = ISIGN(IJ(L),L1(K))
C
C LINES FROM Z(I,J) TO Z(I+1,J) AND Z(I,J+1) ARE NOT SATISFACTORY.
C CONTINUE THE SPIRAL.
C
150 IF (IJ(L).GE.L1(K)) GO TO 170
IJ(L) = IJ(L) + 1
IF (IJ(L).GT.L2(K)) GO TO 160
GO TO 40
160 L2(K) = IJ(L)
IDIR = NXIDIR
GO TO 30
170 IF (IDIR.EQ.NXIDIR) GO TO 180
NXIDIR = NXIDIR + 1
IJ(L) = L1(K)
K = NXIDIR
L = 3 - L
IJ(L) = L2(K)
IF (NXIDIR.GT.3) NXIDIR = 0
GO TO 150
180 IF (IBKEY.NE.0) RETURN
IBKEY = 1
GO TO 10
C
C AN ACCEPTABLE LINE SEGMENT HAS BEEN FOUND.
C FOLLOW THE CONTOUR UNTIL IT EITHER HITS A BOUNDARY OR CLOSES.
C
190 IEDGE = L
CVAL = CV(ICV)
IF (IX.NE.1) IEDGE = IEDGE + 2
IFLAG = 2 + IBKEY

```

GCO 1450
GCO 1460
GCO 1470
GCO 1480
GCO 1490
GCO 1500
GCO 1510
GCO 1520
GCO 1530
GCO 1540
GCO 1550
GCO 1560
GCO 1570
GCO 1580
GCO 1590
GCO 1600
GCO 1610
GCO 1620
GCO 1630
GCO 1640
GCO 1650
GCO 1660
GCO 1670
GCO 1680
GCO 1690
GCO 1700
GCO 1710
GCO 1720
GCO 1730
GCO 1740
GCO 1750
GCO 1760
GCO 1770
GCO 1780
GCO 1790
GCO 1800
GCO 1810
GCO 1820
GCO 1830
GCO 1840
GCO 1850
GCO 1860
GCO 1870
GCO 1880
GCO 1890
GCO 1900
GCO 1910
GCO 1920
GCO 1930
GCO 1940
GCO 1950
GCO 1960
GCO 1970
GCO 1980
GCO 1990
GCO 2000
GCO 2010
GCO 2020
GCO 2030
GCO 2040
GCO 2050
GCO 2060
GCO 2070
GCO 2080
GCO 2090
GCO 2100
GCO 2110
GCO 2120
GCO 2130
GCO 2140
GCO 2150
GCO 2160
GCO 2170
GCO 2180
GCO 2190
GCO 2200


```

XINT(IEDGE) = (CVAL-Z1)/(Z2-Z1) GCO 2210
200 XY(L) = FLOAT(IJ(L)) + XINT(IEDGE) GCO 2220
XY(3-L) = FLOAT(IJ(3-L)) GCO 2230
CALL MARK1(BITMAP, 2*(NX*(NY*(ICV-1)+J-1)+I-1)+L) GCO 2240
CALL DRAW(X, Y, IFLAG+10*ICV) GCO 2250
IF (IFLAG.LT.4) GO TO 210 GCO 2260
ICUR = I GCO 2270
JCUR = J GCO 2280
GO TO 20 GCO 2290
C GCO 2300
C CONTINUE A CONTOUR. THE EDGES ARE NUMBERED CLOCKWISE WITH GCO 2310
C THE BOTTOM EDGE BEING EDGE NUMBER ONE. GCO 2320
C GCO 2330
210 NI = 1 GCO 2340
IF (IEDGE.LT.3) GO TO 220 GCO 2350
I = I - I3(IEDGE) GCO 2360
J = J - I3(IEDGE+2) GCO 2370
220 DO 250 K=1,4 GCO 2380
IF (K.EQ.IEDGE) GO TO 250 GCO 2390
II = I + I3(K) GCO 2400
JJ = J + I3(K+1) GCO 2410
Z1 = Z(II,JJ) GCO 2420
II = I + I3(K+1) GCO 2430
JJ = J + I3(K+2) GCO 2440
Z2 = Z(II,JJ) GCO 2450
IF (CVAL.LE.AMIN1(Z1,Z2)) GO TO 250 GCO 2460
IF (CVAL.GT.AMAX1(Z1,Z2)) GO TO 250 GCO 2470
IF (K.EQ.1) GO TO 230 GCO 2480
IF (K.NE.4) GO TO 240 GCO 2490
230 ZZ = Z1 GCO 2500
Z1 = Z2 GCO 2510
Z2 = ZZ GCO 2520
240 XINT(K) = (CVAL-Z1)/(Z2-Z1) GCO 2530
NI = NI + 1 GCO 2540
KS = K GCO 2550
250 CONTINUE GCO 2560
IF (NI.EQ.2) GO TO 260 GCO 2570
C GCO 2580
C THE CONTOUR CROSSES ALL FOUR EDGES OF THE CELL BEING EXAMINED. GCO 2590
C CHOOSE THE LINES TOP-TO-LEFT AND BOTTOM-TO-RIGHT IF THE GCO 2600
C INTERPOLATION POINT ON THE TOP EDGE IS LESS THAN THE INTERPOLATION GCO 2610
C POINT ON THE BOTTOM EDGE. OTHERWISE, CHOOSE THE OTHER PAIR. THIS GCO 2620
C METHOD PRODUCES THE SAME RESULTS IF THE AXES ARE REVERSED. THE GCO 2630
C CONTOUR MAY CLOSE AT ANY EDGE, BUT MUST NOT CROSS ITSELF INSIDE GCO 2640
C ANY CELL. GCO 2650
C GCO 2660
KS = 5 - IEDGE GCO 2670
IF (XINT(3).LT.XINT(1)) GO TO 260 GCO 2680
KS = 3 - IEDGE GCO 2690
IF (KS.LE.0) KS = KS + 4 GCO 2700
C GCO 2710
C DETERMINE WHETHER THE CONTOUR WILL CLOSE OR RUN INTO A BOUNDARY GCO 2720
C AT EDGE KS OF THE CURRENT CELL. GCO 2730
C GCO 2740
260 L = KS GCO 2750
IFLAG = 1 GCO 2760
ASSIGN 280 TO JUMP GCO 2770
IF (KS.LT.3) GO TO 270 GCO 2780
I = I + I3(KS) GCO 2790
J = J + I3(KS+2) GCO 2800
L = KS - 2 GCO 2810
270 IF (IGET(BITMAP,2*(NX*(NY*(ICV-1)+J-1)+I-1)+L).EQ.0) GO TO 60 GCO 2820
IFLAG = 5 GCO 2830
GO TO 290 GCO 2840
280 IF (IX.NE.0) IFLAG = 4 GCO 2850
290 IEDGE = KS + 2 GCO 2860
IF (IEDGE.GT.4) IEDGE = IEDGE - 4 GCO 2870
XINT(IEDGE) = XINT(KS) GCO 2880
GO TO 200 GCO 2890
C GCO 2900
END GCO 2910

DIMENSION Z(51,51), C(10), WORK(1680) MAN 10
C DIMENSION OF WORK IS LARGE ENOUGH TO CONTAIN MAN 20
C 2*(DIMENSION OF C)*(TOTAL DIMENSION OF Z) USEFUL BITS. SEE THE MAN 30
C BITMAP ROUTINES ACCESSED BY GCONTR. MAN 40

```

REAL MU	MAN	50
EXTERNAL DRAW	MAN	60
COMMON /CUR/ XCUR, YCUR	MAN	70
DATA C(1), C(2), C(3), C(4), C(5) /3.05,3.2,3.5,3.50135,3.6/	MAN	80
DATA C(6), C(7), C(8), C(9), C(10) /3.766413,4.0,4.130149,5.0,	MAN	90
* 10.0/	MAN	100
DATA NX /51/, NY /51/, NF /10/	MAN	110
DATA XMIN /-2.0/, XMAX /2.0/, YMIN /-2.0/, YMAX /2.0/, MU /0.3/	MAN	120
DX = (XMAX-XMIN)/FLOAT(NX-1)	MAN	130
DY = (YMAX-YMIN)/FLOAT(NY-1)	MAN	140
XCUR = 1.0	MAN	150
YCUR = 1.0	MAN	160
IF (MOD(NX,2).NE.0) YCUR = FLOAT(NY)	MAN	170
IF (MOD(NY,2).NE.0) XCUR = FLOAT(NX)	MAN	180
X = XMIN - DX	MAN	190
DO 20 I=1,NX	MAN	200
Y = YMIN - DY	MAN	210
X = X + DX	MAN	220
DO 10 J=1,NY	MAN	230
Y = Y + DY	MAN	240
Z(I,J) = (1.0-MU)*(2.0/SQRT((X-MU)**2+Y**2)+(X-MU)**2+Y**2)	MAN	250
* + MU*(2.0/SQRT((X+1.0-MU)**2+Y**2)+(X+1.0-MU)**2+Y**2)	MAN	260
10 CONTINUE	MAN	270
20 CONTINUE	MAN	280
CALL GCONTR(Z, 51, NX, NY, C, NF, 1.E6, WORK, DRAW)	MAN	290
STOP	MAN	300
END	MAN	310
REAL Z(51,51), C(10), CVAL(10), MU	MAN	10
INTEGER WORK(1680), L(10), CLAB(10)	MAN	20
C DIMENSION OF WORK IS LARGE ENOUGH TO CONTAIN	MAN	30
C 2*(DIMENSION OF C)*(TOTAL DIMENSION OF Z) USEFUL BITS. SEE THE	MAN	40
C BITMAP ROUTINES ACCESSED BY GCONTR.	MAN	50
EXTERNAL DRAW	MAN	60
COMMON /GCTCOM/ XCUR, YCUR, XL, YL, CVAL, CLAB, NCH	MAN	70
DATA C(1), C(2), C(3), C(4), C(5) /3.05,3.2,3.5,3.50135,3.6/	MAN	80
DATA C(6), C(7), C(8), C(9), C(10) /3.766413,4.0,4.130149,5.0,	MAN	90
* 10.0/	MAN	100
DATA L(1), L(2), L(3), L(4), L(5) /1HA,1HB,1HC,1HD,1HE/	MAN	110
DATA L(6), L(7), L(8), L(9), L(10) /1HF,1HG,1HH,1HI,1HJ/	MAN	120
DATA NX /51/, NY /51/, NF /10/, NXG /5/, NYG /5/	MAN	130
DATA XMIN /-2.0/, XMAX /2.0/, YMIN /-2.0/, YMAX /2.0/, MU /0.3/	MAN	140
DATA XLEN /8.0/, YLEN /8.0/	MAN	150
C INITIALIZE PLOTTING SUBROUTINES.	MAN	160
CALL PLOTS	MAN	170
DX = (XMAX-XMIN)/FLOAT(NX-1)	MAN	180
DY = (YMAX-YMIN)/FLOAT(NY-1)	MAN	190
XL = XLEN/FLOAT(NX)	MAN	200
YL = YLEN/FLOAT(NY)	MAN	210
XCUR = 1.0	MAN	220
YCUR = 1.0	MAN	230
IF (MOD(NX,2).NE.0) YCUR = FLOAT(NY)	MAN	240
IF (MOD(NY,2).NE.0) XCUR = FLOAT(NX)	MAN	250
X = XMIN - DX	MAN	260
DO 20 I=1,NX	MAN	270
Y = YMIN - DY	MAN	280
X = X + DX	MAN	290
DO 10 J=1,NY	MAN	300
Y = Y + DY	MAN	310
C EVALUATE FUNCTION TO BE PLOTTED.	MAN	320
Z(I,J) = (1.0-MU)*(2.0/SQRT((X-MU)**2+Y**2)+(X-MU)**2+Y**2)	MAN	330
* + MU*(2.0/SQRT((X+1.0-MU)**2+Y**2)+(X+1.0-MU)**2+Y**2)	MAN	340
10 CONTINUE	MAN	350
20 CONTINUE	MAN	360
DO 30 I=1,NF	MAN	370
CVAL(I) = C(I)	MAN	380
CLAB(I) = L(I)	MAN	390
30 CONTINUE	MAN	400
NCH = 1	MAN	410
C PEN UP MOVE TO BELOW LOWER LEFT CORNER OF PAGE.	MAN	420
C THIS CALL WORKS DIFFERENTLY ON DIFFERENT MACHINES. YOU MAY	MAN	430
C NEED TO CHANGE IT.	MAN	440
CALL PLOT(0.0, -11.0, -3)	MAN	450
C PEN UP MOVE TO 1 INCH ABOVE LOWER LEFT CORNER OF PAGE.	MAN	460
CALL PLOT(0.0, 1.0, -3)	MAN	470

	SX = 8.0/FLOAT(NXG)	MAN	480
	SY = 8.0/FLOAT(NXG)	MAN	490
C	DRAW A GRID.	MAN	500
	CALL CGRID(1, NXG, SX, 0.0, 0.0, NYG, SY, 0.0, 0.0)	MAN	510
C	DRAW THE CONTOUR PLOTS.	MAN	520
	CALL GCONTR(Z, 51, NX, NY, CVAL, NF, 1.0E6, WORK, DRAW)	MAN	530
	XX = 9.0	MAN	540
	YY = 8.0	MAN	550
C	WRITE A TABLE OF CONTOUR LABELS AND VALUES.	MAN	560
	CALL SYMBOL(XX, YY+0.14, 0.07, 10HCONTOUR ID, 0.0, 10)	MAN	570
	DO 40 I=1,NF	MAN	580
	CALL SYMBOL(XX, YY, 0.07, L(I), 0.0, 2)	MAN	590
	CALL NUMBER(XX+0.12, YY, 0.07, C(I), 0.0, 5)	MAN	600
	YY = YY - 0.14	MAN	610
	40 CONTINUE	MAN	620
C	PEN UP MOVE TO BELOW LOWER RIGHT CORNER OF PAGE.	MAN	630
C	THIS CALL WORKS DIFFERENTLY ON DIFFERENT MACHINES. YOU MAY NEED	MAN	640
C	TO CHANGE IT, OR YOU MAY NOT NEED IT.	MAN	650
	CALL PLOT(10.0, -11.0, -3)	MAN	660
C	REDUCE PICTURE SIZE, PLOT END OF FILE INFORMATION.	MAN	670
C	THE END OF FILE INFORMATION MAY NOT BE AVAILABLE AT ALL SITES.	MAN	680
C	IF NOT AVAILABLE, CHANGE THE NEXT TWO STATEMENTS TO COMMENTS.	MAN	690
	CALL FACTOR(0.3)	MAN	700
	CALL PLOT(0.0, 0.0, 999)	MAN	710
	STOP	MAN	720
C		MAN	730
	END	MAN	740
	SUBROUTINE DRAW(X, Y, IFLAG)	DRA	10
C	THIS SUBROUTINE USES CALCOMP PLOT ROUTINES TO DRAW LINES FOR THE	DRA	20
C	CONTOUR PLOTTING ROUTINE GCONTR.	DRA	30
	REAL CVAL(10)	DRA	40
	INTEGER CLAB(10)	DRA	50
	COMMON /GCTCOM/ XCUR, YCUR, XL, YL, CVAL, CLAB, NCH	DRA	60
	DATA IBLANK /1H /	DRA	70
	IH = IFLAG/10	DRA	80
	IL = IFLAG - 10*IH	DRA	90
	IF (IL.EQ.6) GO TO 40	DRA	100
	IPEN = 2	DRA	110
	IF (IL.EQ.2) IPEN = 3	DRA	120
	IF (IL.EQ.3) IPEN = 3	DRA	130
	XCUR = X	DRA	140
	YCUR = Y	DRA	150
	XX = (X-1.0)*XL	DRA	160
	YY = (Y-1.0)*YL	DRA	170
	CALL PLOT(XX, YY, IPEN)	DRA	180
	IF (IL.LT.2) GO TO 30	DRA	190
	IF (IL.GT.4) GO TO 30	DRA	200
	IF (NCH.LT.1) GO TO 30	DRA	210
	IF (CLAB(IH).EQ.IBLANK) GO TO 30	DRA	220
	IF (CLAB(IH).NE.0) GO TO 10	DRA	230
	CALL NUMBER(XX, YY-0.03, 0.07, CVAL(IH), 0.0, -1)	DRA	240
	GO TO 20	DRA	250
	10 CALL SYMBOL(XX, YY-0.03, 0.07, CLAB(IH), 0.0, NCH)	DRA	260
	20 CALL PLOT(XX, YY, 3)	DRA	270
	30 RETURN	DRA	280
	40 X = XCUR	DRA	290
	Y = YCUR	DRA	300
	RETURN	DRA	310
C		DRA	320
	END	DRA	330
	SUBROUTINE DRAW(X, Y, IFLAG)	DRA	10
C		DRA	20
C	DO OUTPUT FOR GCONTR.	DRA	30
C		DRA	40
	INTEGER PRINT	DRA	50
	COMMON /CUR/ XCUR, YCUR	DRA	60
	DATA PRINT /6/	DRA	70
C	PRINT IS THE SYSTEM PRINTER FORTRAN I/O UNIT NUMBER.	DRA	80
	ICONT = IFLAG/10	DRA	90
	JUMP = MOD(IFLAG,10)	DRA	100
	GO TO (10, 20, 30, 40, 50, 60), JUMP	DRA	110
	10 WRITE (PRINT,99999) ICONT, X, Y	DRA	120
	GO TO 70	DRA	130

20	WRITE (PRINT,99998) ICONT, X, Y	DRA	140
	GO TO 70	DRA	150
30	WRITE (PRINT,99997) ICONT, X, Y	DRA	160
	GO TO 70	DRA	170
40	WRITE (PRINT,99996) ICONT, X, Y	DRA	180
	GO TO 70	DRA	190
50	WRITE (PRINT,99995) ICONT, X, Y	DRA	200
	GO TO 70	DRA	210
60	WRITE (PRINT,99994)	DRA	220
	X = XCUR	DRA	230
	Y = YCUR	DRA	240
70	RETURN	DRA	250
99999	FORMAT (17H CONTINUE CONTOUR, I3, 3H TO, 1P2E14.7)	DRA	260
99998	FORMAT (14H START CONTOUR, I3, 19H ON THE BOUNDARY AT, 1P2E14.7)	DRA	270
99997	FORMAT (14H START CONTOUR, I3, 19H IN THE INTERIOR AT, 1P2E14.7)	DRA	280
99996	FORMAT (15H FINISH CONTOUR, I3, 19H ON THE BOUNDARY AT, 1P2E14.7)	DRA	290
99995	FORMAT (15H FINISH CONTOUR, I3, 19H IN THE INTERIOR AT, 1P2E14.7)	DRA	300
99994	FORMAT (33H REQUEST FOR CURRENT PEN POSITION)	DRA	310
	END	DRA	320
	SUBROUTINE CGRID(NOPT, NX, SX, XS, XF, NY, SY, YS, YF)	CGR	10
C		CGR	20
C	SUBROUTINE WHICH DRAWS A FRAME AROUND THE PLOT AND DRAWS	CGR	30
C	EITHER TICK MARKS OR GRID LINES.	CGR	40
C		CGR	50
C	PARAMETERS: NOPT -- =0, DRAW TICKS ONLY	CGR	60
C	=1, DRAW GRID LINES	CGR	70
C	=2, DRAW GRID LINES TO EDGE OF FRAME.	CGR	80
C	NX -- NUMBER OF INTERVALS IN X DIRECTION	CGR	90
C	SX -- SPACING IN INCHES BETWEEN TICK MARKS OR GRID LINES	CGR	100
C	ALONG THE X AXIS	CGR	110
C	XS -- LOCATION OF FIRST TICK OR GRID LINE ON X AXIS	CGR	120
C	XF -- LOCATION OF RIGHT EDGE OF FRAME	CGR	130
C	NY -- NUMBER OF INTERVALS IN Y DIRECTION	CGR	140
C	SY -- SPACING IN INCHES BETWEEN TICK MARKS OR GRID LINES	CGR	150
C	ALONG THE Y AXIS	CGR	160
C	YS -- LOCATION OF FIRST TICK OR GRID LINE ON Y AXIS	CGR	170
C	YF -- LOCATION OF TOP EDGE OF FRAME	CGR	180
C	ASSUMPTIONS: NX, SX, NY, SY ALL POSITIVE.	CGR	190
C	THE LOWER LEFT-HAND CORNER OF THE FRAME IS DRAWN AT (0,0)	CGR	200
C	IF XS10, USE 0; IF YS10, USE 0	CGR	210
C	IF XF1=0, USE NX*SX; IF YF1=0, USE NY*SY.	CGR	220
C		CGR	230
	XINC = SX	CGR	240
	YINC = SY	CGR	250
	XLGTH = FLOAT(NX)*SX	CGR	260
	YLGTH = FLOAT(NY)*SY	CGR	270
	XMIN = AMAX1(XS,0.0)	CGR	280
	YMIN = AMAX1(YS,0.0)	CGR	290
	XMAX = AMAX1(XF,XLGTH+XMIN)	CGR	300
	YMAX = AMAX1(YF,YLGTH+YMIN)	CGR	310
C		CGR	320
C	DRAW FRAME.	CGR	330
C		CGR	340
	CALL PLOT(0.0, 0.0, 3)	CGR	350
	CALL PLOT(XMAX, 0.0, 2)	CGR	360
	CALL PLOT(XMAX, YMAX, 2)	CGR	370
	CALL PLOT(0.0, YMAX, 2)	CGR	380
	CALL PLOT(0.0, 0.0, 2)	CGR	390
	IF (NOPT.NE.0) GO TO 130	CGR	400
C		CGR	410
C	DRAW TICK MARKS.	CGR	420
C		CGR	430
	DO 120 J=1,4	CGR	440
	GO TO (10, 50, 20, 40), J	CGR	450
10	X2 = 0.0	CGR	460
	IF (XMIN.NE.0.0) X2 = XMIN - SX	CGR	470
	Y2 = 0.0	CGR	480
	GO TO 30	CGR	490
20	XINC = -SX	CGR	500
	X2 = XMIN + XLGTH + SX	CGR	510
	IF (XMAX.EQ.XMIN+XLGTH) X2 = XMAX	CGR	520
	Y2 = YMAX	CGR	530
30	Y1 = Y2	CGR	540
	Y2 = Y2 + SIGN(0.125,XINC)	CGR	550

	N = NX	CGR 560
	IF (ABS(XMAX-XMIN-XLGTH)+ABS(XMIN)) 70, 80, 70	CGR 570
40	YINC = -SY	CGR 580
	Y2 = YMIN + YLGTH + SY	CGR 590
	IF (YMAX.EQ.YMIN+YLGTH) Y2 = YMAX	CGR 600
	X2 = 0.0	CGR 610
	GO TO 60	CGR 620
50	Y2 = 0.0	CGR 630
	IF (YMIN.NE.0.0) Y2 = YMIN - SY	CGR 640
	X2 = XMAX	CGR 650
60	X1 = X2	CGR 660
	N = NY	CGR 670
	X2 = X2 - SIGN(0.125,YINC)	CGR 680
	IF (ABS(YMAX-YMIN-YLGTH)+ABS(YMIN)) 70, 80, 70	CGR 690
70	N = N + 1	CGR 700
80	DO 110 I=1,N	CGR 710
	IF (MOD(J,2).EQ.0) GO TO 90	CGR 720
	X2 = X2 + XINC	CGR 730
	X1 = X2	CGR 740
	GO TO 100	CGR 750
90	Y2 = Y2 + YINC	CGR 760
	Y1 = Y2	CGR 770
100	CALL PLOT(X1, Y1, 3)	CGR 780
	CALL PLOT(X2, Y2, 2)	CGR 790
110	CONTINUE	CGR 800
120	CONTINUE	CGR 810
	GO TO 240	CGR 820
C		CGR 830
C	DRAW GRID LINES	CGR 840
C		CGR 850
130	X1 = XMIN	CGR 860
	X2 = XMIN + XLGTH	CGR 870
	IF (NOPT.NE.2) GO TO 140	CGR 880
	X1 = 0.0	CGR 890
	X2 = XMAX	CGR 900
140	Y1 = YMIN - SY	CGR 910
	N = NY + 1	CGR 920
	IF (YMAX.EQ.YMIN+YLGTH) N = N - 1	CGR 930
	IF (YMIN.NE.0.0) GO TO 150	CGR 940
	Y1 = 0.0	CGR 950
	N = N - 1	CGR 960
150	IF (N.LE.0) GO TO 170	CGR 970
	J = 1	CGR 980
	DO 160 I=1,N	CGR 990
	J = -J	CGR 1000
	Y1 = Y1 + SY	CGR 1010
	CALL PLOT(X1, Y1, 3)	CGR 1020
	CALL PLOT(X2, Y1, 2)	CGR 1030
	XX = X1	CGR 1040
	X1 = X2	CGR 1050
	X2 = XX	CGR 1060
160	CONTINUE	CGR 1070
170	Y1 = YMIN + YLGTH	CGR 1080
	Y2 = YMIN	CGR 1090
	IF (NOPT.NE.2) GO TO 180	CGR 1100
	Y1 = YMAX	CGR 1110
	Y2 = 0.0	CGR 1120
180	N = NX + 1	CGR 1130
	IF (J.LT.0) GO TO 200	CGR 1140
	X1 = XMIN - SX	CGR 1150
	IF (XMAX.EQ.XMIN+XLGTH) N = N - 1	CGR 1160
	IF (XMIN.NE.0.0) GO TO 190	CGR 1170
	X1 = 0.0	CGR 1180
	N = N - 1	CGR 1190
190	IF (N.LE.0) GO TO 240	CGR 1200
	XINC = SX	CGR 1210
	GO TO 220	CGR 1220
200	X1 = XMIN + XLGTH + SX	CGR 1230
	IF (XMIN.EQ.0.0) N = N - 1	CGR 1240
	IF (XMAX.NE.XLGTH+XMIN) GO TO 210	CGR 1250
	N = N - 1	CGR 1260
	X1 = XMAX	CGR 1270
210	XINC = -SX	CGR 1280
220	DO 230 I=1,N	CGR 1290
	X1 = X1 + XINC	CGR 1300
	CALL PLOT(X1, Y1, 3)	CGR 1310

COLLECTED ALGORITHMS (cont.)

531-P12- 0

```
      CALL PLOT(X1, Y2, 2)
      XX = Y1
      Y1 = Y2
      Y2 = XX
230 CONTINUE
240 RETURN
C
      END
```

```
CGR 1320
CGR 1330
CGR 1340
CGR 1350
CGR 1360
CGR 1370
CGR 1380
CGR 1390
```

ALGORITHM 532

Software for Roundoff Analysis [Z]

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and

DAVID SPOONER

Pennsylvania State University

Key Words and Phrases: automatic roundoff analysis, numerical stability, numerical linear algebra

CR Categories: 5.10, 5.11, 5.14

Language: Fortran

DESCRIPTION

This software package is a complement to [1] where its usage and performance are described.

REFERENCES

1. MILLER, W., AND SPOONER, D. Software for roundoff analysis, II. *ACM Trans. Math. Software* 4, 4 (Dec. 1978), 369-387.

ALGORITHM

NAME(n): indicates a Fortran module with n records

NAME^D: indicates test data for the minicompiler

Contents for the minicompiler: DATA(570), MAIN(388), CODGEN(445), NEXT(12), CODOPT(148), INSERT(69), GETNAM(32), GETDIM(34), RDIM(51), ADD(114), STORE(69), GETVAL(88), SYMINT(39), LEXAN(422), GETSTM(343), K FIND(118), INTERP(322), REALOP(25), OPER(504), FINISH(203), ADDTMP(70), ADDSUB(51), ADDNAM(47), TERROR(58), ERROR(80), IERROR(95)

```

C          *****
C          *                               *
C          *      OVERVIEW      *
C          *                               *
C          *****
C
C
C          THIS PROGRAM IS A COMPILER WHICH TAKES AS DATA A PROGRAM WRITTEN
C          IN A SIMPLE PROGRAMMING LANGUAGE ( A DESCRIPTION OF WHICH FOLLOWS)
C          AND PRODUCES AS OUTPUT A TRANSLATION OF THAT PROGRAM INTO A SERIES
C          OF ASSIGNMENT STATEMENTS.  THE STRAIGHT LINE CODE TRANSLATION IS

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C PRINTED IN A READABLE FORM FOR THE COMPILER USER AND ALSO PUNCHED
 C IN A CODED FORM FOR USE AS INPUT TO ANY OF SEVERAL ROUNDOFF ERROR
 C ANALYSIS PROGRAMS OF WEBB MILLER.
 C THE MAIN ROUTINE OF THE COMPILER IS AN LL-1 PARSING ROUTINE. IT
 C CALLS THE LEXICAL ANALYZER TO PRODUCE INTERNAL TOKENS BY SCANNING
 C THE SOURCE CODE. THE PARSER ALSO CALLS THE ROUTINE CODGEN WHICH
 C PRODUCES INTERMEDIATE CODE FOR THE INTERPRETER AND WHICH PERFORMS
 C SOME SYNTAX CHECKING AND OTHER PERIPHERAL FUNCTIONS NECESSARY
 C FOR PARSING.
 C WHEN THE ENTIRE INPUT HAS BEEN SCANNED AND PARSED, THE INTERPRETING
 C ROUTINES ARE ACTIVATED. THESE INTERPRET THE INTERMEDIATE CODE AND
 C GENERATE THE FINAL STRAIGHT LINE PRINTED AND PUNCHED OUTPUT. DURING
 C INTERPRETATION, ALL INTEGER EXPRESSIONS ARE ACTUALLY EVALUATED IN
 C ORDER TO PERFORM THE CORRECT NUMBER OF ITERATIONS OF EXPLICIT FOR-
 C LOOPS AND OF FOR-LOOPS IMPLICIT IN SUMMATION EXPRESSIONS, AND TO
 C INTERPRETIVELY PERFORM IF-THEN TESTS. IN CONTRAST, NO ACTUAL
 C REAL ARITHMETIC COMPUTATION IS DONE. THROUGHOUT, ALL REAL
 C VARIABLES ARE TREATED SYMBOLICALLY AS BEING THE N-TH INPUT VALUE,
 C INTERMEDIATE VALUE, OR REAL CONSTANT.
 C BOTH THE LEXICAL ANALYZER AND THE INTERPRETER COMPRISE SEVERAL SUB-
 C ROUTINES. IN ADDITION THERE ARE COLLECTIONS OF SYMBOL TABLE ROUTINES,
 C ERROR ROUTINES AND ROUTINES TO PRINT AND PUNCH THE COMPILER'S
 C OUTPUT.
 C IF ANY ERRORS ARE ENCOUNTERED DURING PARSING, THE PARSER CONTINUES
 C TO CHECK FOR SYNTAX ERRORS, BUT NO FURTHER INTERMEDIATE CODE WILL BE
 C GENERATED, AND NO INTERPRETATION WILL TAKE PLACE. SIMILARLY, IF
 C EXECUTION ERRORS ARE IDENTIFIED DURING INTERPRETATION, NO STRAIGHT
 C LINE CODE WILL BE EITHER PRINTED OR PUNCHED.

```

*****
*                               *
*           THE SOURCE LANGUAGE   *
*                               *
*****

```

C THE LANGUAGE TO BE COMPILED IS A SIMPLE LANGUAGE DESIGNED FOR
 C CODING NUMERICAL ALGORITHMS. IT BASICALLY INCLUDES REAL ASSIGN-
 C MENT STATEMENTS, DIMENSION STATEMENTS AND SOME BLOCK STRUCTURE
 C IMPOSED BY FOR-LOOPS, AND IF-THEN TESTS. THERE ARE NO MIXED-MODE
 C ARITHMETIC EXPRESSIONS, AND NO STATEMENT LABELS. INTEGER EXPRESSIONS
 C AND VARIABLES ARE USED ONLY FOR DIMENSIONING REAL ARRAYS, FOR
 C DEFINING BOUNDS IN FOR-LOOPS AND SUMMATION EXPRESSION LOOPS, AND FOR
 C VARIABLES TO BE TESTED IN IF-THEN STATEMENTS.

C THE CARD FORMAT IS SIMILAR TO FORTRAN. STATEMENTS APPEAR
 C IN COLUMNS 7-72, AND A 1 IN COLUMN 6 INDICATES CONTINUATION. A
 C C IN COLUMN 1 INDICATES A COMMENT.

C THERE ARE SEVEN STATEMENT TYPES, BRIEFLY DESCRIBED BELOW:

C 1. THE TEST STATEMENT

C USE: TO ASSIGN VALUES TO INTEGER VARIABLES WHICH WILL THEMSELVES
 C BE USED TO DIMENSION REAL ARRAYS.

C FORM: TEST(I1,I2,I3,...)

C WHERE EACH I IS AN ASSIGNMENT STATEMENT OF THE FORM
 C INTEGER VARIABLE = INTEGER CONSTANT.

C 2. THE DIMENSION STATEMENT

C USE: TO ASSIGN DIMENSIONS TO ARRAYS. THE ARRAY NAME MAY BE
 C EITHER A DEFAULT REAL OR INTEGER IDENTIFIER. ITS USE IN
 C THE DIMENSION STATEMENT CONSTITUTES AN IMPLICIT REAL
 C DECLARATION.

C FORM: DIMENSION(D1,D2,D3 ..)

C WHERE EACH D IS OF ONE OF THE FOLLOWING FORMS:

C IDENTIFIER(I)

C IDENTIFIER(I1,I2)

C AND EACH I IS AN INTEGER VARIABLE OR INTEGER CONSTANT.

C NOTE: DIMENSION AND TEST STATEMENTS ARE NON-EXECUTABLE AND MUST
 C PRECEED ALL EXECUTABLE STATEMENTS. IN ADDITION, AN INTEGER

00000600
 00000650
 00000700
 00000750
 00000800
 00000850
 00000900
 00000950
 00001000
 00001050
 00001100
 00001150
 00001200
 00001250
 00001300
 00001350
 00001400
 00001450
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 00003900
 00003950
 00004000
 00004050
 00004100
 00004150
 00004200
 00004250
 00004300
 00004350

C VARIABLE APPEARING IN A DIMENSION STATEMENT MUST BE ASSIGNED A 00004400
C VALUE IN A PRECEDING TEST STATEMENT. 00004450
C 00004500
C 3. THE INPUT/OUTPUT STATEMENTS 00004550
C 00004600
C USE: TO NOTIFY THE COMPILER THAT CERTAIN VALUES WILL BE SUPPLIED 00004650
C BY THE PROGRAMMER AS INITIAL VALUES FOR REAL VARIABLES WHEN 00004700
C THE STRAIGHT LINE CODE IS USED AS INPUT FOR A ROUND OFF ERROR 00004750
C ANALYSIS; OR THAT CERTAIN REAL VARIABLES ARE EXPECTED TO 00004800
C RECEIVE VALUES AS A RESULT OF RUNNING THE PROGRAM BEING 00004850
C COMPILED. 00004900
C 00004950
C FORM: INPUT(D1,D2,D3,...) 00005000
C OUTPUT(D1,D2,D3,...) 00005050
C 00005100
C EACH D IS EITHER THE NAME OF A REAL SCALAR, A SINGLE ARRAY 00005150
C ELEMENT, OR AN ENTIRE ARRAY. IN THE LATTER CASE THE ARRAY 00005200
C WILL BE INPUT(OUTPUT) IN COLUMN MAJOR ORDER. 00005250
C NOTE THAT ONLY VALUES WHICH ARE THE RESULT OF A 00005300
C COMPUTATION MAY BE OUTPUT. CONSTANTS OR DATA VALUES 00005350
C MAY NOT BE OUTPUT. 00005400
C 00005450
C 4. THE REAL ASSIGNMENT STATEMENT 00005500
C USE: TO ASSIGN A VALUE TO A REAL VARIABLE 00005600
C 00005650
C FORM: REAL VARIABLE = REAL EXPRESSION 00005700
C 00005750
C WHERE THE REAL VARIABLE IS EITHER A REAL SCALAR OR SINGLE 00005800
C ARRAY ELEMENT. REAL EXPRESSIONS ARE MADE UP OF REAL 00005850
C VARIABLES AND CONSTANTS COMBINED WITH THE BINARY OPERATORS 00005900
C +, -, * AND / AND THE UNARY OPERATORS UNARY - AND SQRT. 00005950
C (THE OPERAND OF THE SQRT MUST APPEAR IN PARENTHESES.) 00006000
C (NOTE THAT THERE IS NO REAL EXPONENTIATION ALLOWED.) 00006050
C OPERATOR PRECEDENCE IS AS IN STANDARD FORTRAN. IN ADDITION 00006100
C THERE IS A SUMMATION OPERATION ON ARRAY VECTORS, IN EFFECT, 00006150
C A BUILT-IN INNER PRODUCT. A SUMMATION EXPRESSION CAN APPEAR 00006200
C IN A REAL EXPRESSION AND IS OF THE FORM: 00006250
C 00006300
C SUMMATION(D1 * D2, SUMMATION-VARIABLE = INTEXP1 TO INTEXP2) 00006350
C 00006400
C WHERE INTEXP IS AN INTEGER EXPRESSION, THE SUMMATION 00006450
C VARIABLE IS ANY INTEGER IDENTIFIER NAME, AND WHERE D1 AND D2 00006500
C ARE EACH OF ONE OF THE FORMS: 00006550
C ARRAY NAME(SUMMATION-VARIABLE) 00006600
C ARRAY NAME(SUMMATION-VARIABLE, SUMMATION-VARIABLE) 00006650
C ARRAY NAME(SUMMATION-VARIABLE, INTEXP) 00006700
C ARRAY NAME(INTEXP, SUMMATION-VARIABLE) 00006750
C 00006800
C A SUMMATION EXPRESSION WILL BE INTERPRETED AS AN IMPLICIT 00006850
C FOR-LOOP. ANY USE OF THE SUMMATION VARIABLE IN THE INTEGER 00006900
C EXPRESSIONS BOUNDING THE SUMMATION LOOP WILL BE FLAGGED AS 00006950
C AN ERROR. 00007000
C 00007050
C 5. THE FOR STATEMENT 00007100
C 00007150
C USE: AS A MEANS OF INDICATING THAT A BLOCK OF STATEMENTS IS TO BE 00007200
C ITERATIVELY EXECUTED A SPECIFIED NUMBER OF TIMES. 00007250
C 00007300
C FORM: FOR INTEGER-VARIABLE = INTEXP1 TO INTEXP2 BY INCREMENT 00007350
C 00007400
C WHERE INTEXP STANDS FOR INTEGER EXPRESSION, AND INCREMENT IS 00007450
C WRITTEN AS EITHER +1, -1 OR 1. 00007500
C 00007550
C INTERPRETATION: ALL STATEMENTS UP TO THE END STATEMENT MATCHING 00007600
C THIS FOR STATEMENT (SEE BELOW FOR DISCUSSION OF END 00007650
C STATEMENTS) WILL BE ITERATIVELY INTERPRETED AS IN A FORTRAN 00007700
C DO LOOP, EXCEPT THAT NEGATIVE INCREMENTS ARE ALLOWED AND IN 00007750
C THIS CASE, THE LOOP VARIABLE TEST IS DONE AT THE TOP OF THE 00007800
C LOOP. THUS, EMPTY LOOPS ARE POSSIBLE, THAT IS THOSE WHICH 00007850
C WILL NOT BE EXECUTED AT ALL. (NOTE THE SAME IS TRUE OF THE 00007900
C IMPLICIT FOR-LOOP IN A SUMMATION EXPRESSION.) 00007950
C 00008000
C 6. THE IF-THEN STATEMENT 00008050
C 00008100
C USE: TO ALLOW SELECTIVE EXECUTION OF A BLOCK OF STATEMENTS 00008150

C DEPENDING ON THE OUTCOME OF A COMPARISON OF THE VALUES OF 00008200
C TWO INTEGER EXPRESSIONS. 00008250
C 00008300
C FORM: IF INTEXP1 .R. INTEXP2 THEN 00008350
C 00008400
C WHERE INTEXP STANDS FOR INTEGER EXPRESSION AND R IS ONE OF 00008450
C THE RELATIONS WRITTEN EQ,NE,GT,LT,LE OR GE WITH THE 00008500
C STANDARD FORTRAN DENOTATION. NOTE THAT THERE ARE NO 00008550
C PARENTHESES AROUND THE RELATIONAL EXPRESSION. 00008600
C 00008650
C INTERPRETATION: IF THE TEST SUCCEEDS, THAT IS IF THE TWO INTEGER 00008700
C EXPRESSIONS ARE RELATED IN THE INDICATED WAY AT THE TIME OF 00008750
C INTERPRETATION, THEN ALL THE STATEMENTS UP TO THE NEXT END 00008800
C STATEMENT WILL BE INTERPRETED. OTHERWISE, THE FIRST 00008850
C EXECUTABLE STATEMENT FOLLOWING THE NEXT END STATEMENT WILL 00008900
C BE THE NEXT STATEMENT INTERPRETED. 00008950
C 00009000
C 7. THE END STATEMENT 00009050
C 00009100
C USE: TO DEFINE THE ENDS OF BLOCKS OF STATEMENTS BEGINNING WITH 00009150
C FOR STATEMENTS OR IF-THEN STATEMENTS. 00009200
C 00009250
C FORM1: END 00009300
C FORM2 : END(INTEGER-VARIABLE) 00009350
C 00009400
C MEANING: WHEN FORM 1 IS USED THE EFFECT IS TO CLOSE THE BLOCK OF 00009450
C STATEMENTS BEGINNING AT THE NEAREST PRECEDING FOR OR 00009500
C IF-THEN STATEMENT. 00009550
C WHEN FORM2 IS USED THE EFFECT IS TO CLOSE THE FOR BLOCK 00009600
C WHOSE LOOP VARIABLE MATCHES THE END STATEMENT VARIABLE. IN 00009650
C ADDITION, ANY FOR OR IF-THEN BLOCKS WHICH BEGIN BETWEEN THIS 00009700
C END STATEMENT AND ITS MATCHING FOR STATEMENT ARE CLOSED. 00009750
C THIS INTERPRETATION IMPOSES STANDARD FORTRAN LIKE NESTING 00009800
C CONVENTIONS ON FOR AND IF-THEN BLOCKS. THAT IS, A SEQUENCE 00009850
C OF STATEMENTS 00009900
C FOR K = IE1 TO IE2 BY 1 00009950
C . 00010000
C . 00010050
C FOR I = IE3 TO IE4 BY 1 00010100
C . 00010150
C . 00010200
C END(K) 00010250
C . 00010300
C END(I) 00010350
C WILL RESULT IN AN ERROR MESSAGE WHEN THE END(I) STATEMENT 00010400
C IS ENCOUNTERED, BECAUSE BOTH FOR STATEMENTS WILL HAVE BEEN 00010450
C CLOSED BY THE PARSER WHEN THE END(K) STATEMENT WAS PARSED. 00010500
C 00010550
C NOTE: ADDITIONAL RESTRICTIONS ON BLOCK STRUCTURES: 00010600
C 1) AT MOST EIGHT FOR AND/OR IF-THEN BLOCKS CAN BE BEGUN 00010650
C BEFORE AN END STATEMENT OCCURS. 00010700
C 2) A FOR LOOP VARIABLE CANNOT BE USED AGAIN AS AN EXPLICIT 00010750
C FOR LOOP VARIABLE WITHIN ITS ORIGINAL LOOP. IT CAN BE 00010800
C REUSED AS A SUMMATION VARIABLE, HOWEVER. 00010850
C 00010900
C 00010950
C 8. THE STOP STATEMENT 00011000
C 00011050
C USE: TO DENOTE THE END OF THE PROGRAM 00011100
C 00011150
C FORM: *STOP 00011200
C 00011250

Contents of program comparing rounding error in a single algorithm with perturbations of the problem: MAIN1(372), F(64), ROUND(116), GETRHO(64), OMEGA(76)

C 00000050
C THE USER SUPPLIES: 00000100
C 00000150
C THE OUTPUT FROM THE MINICOMPILER. 00000200
C 00000250
C THE ENTRIES OF THE INITIAL SET OF DATA. ONE ENTRY PER CARD, EACH 00000300
C ENTRY WRITTEN WITH A DECIMAL POINT AND CONTAINED IN THE FIRST 00000350
C TWENTY COLUMNS. 00000400
C 00000450
C THE CHOICE OF OMEGA (+) OR RHO (-). FORMAT(I2). 00000500

```

C      CODE : PERTURBATIONS                                00000550
C      Ø1 : IN THE DATA, MEASURED COMPONENT-WISE.        00000600
C      Ø2 : IN THE DATA AND RESULT, MEASURED COMPONENT-WISE. 00000650
C      Ø3 : IN THE DATA, MEASURED NORM-WISE.             00000700
C      Ø4 : IN THE DATA AND RESULT, MEASURED NORM-WISE.   00000750
C
C      THE STOPPING VALUE FOR THE MAXIMIZER. WRITTEN WITH A DECIMAL
C      POINT AND CONTAINED IN THE FIRST TWENTY COLUMNS OF ITS CARD. 00000850
C
C THE SOFTWARE RETURNS:                                    00000900
C
C      AN ANNOTATED LISTING OF THE USER-SUPPLIED INFORMATION, PLUS THE
C      ERROR-COMPARING VALUE, THE CONSTRAINT VALUES (IF ANY) AND
C      THE OUTPUT COMPUTED AT THE INITIAL SET OF DATA.     00000950
C
C      A LIST OF SELECTED VALUES FOUND BY THE MAXIMIZER.   00001000
C
C      THE FINAL SET OF DATA.                              00001050
C
C      IF INSTABILITY IS DIAGNOSED, THEN ALL ARITHMETIC OPERATIONS
C      AT THE FINAL SET OF DATA ARE LISTED.               00001100
C
C OTHER INFORMATION IS RETURNED IF EXCEPTIONS ARISE.     00001150
C
C ----- 00001200
C
C THE USER CAN AVOID THE MINICOMPILER BY SUPPLYING:     00001250
C
C      THE NUMBER OF OPERATIONS IN THE PROGRAM BEING TESTED. GE.1
C      AND LE.2ØØ. FORMAT(13). (INSTRUCTIONS ARE PROVIDED BELOW FOR
C      RAISING THE UPPER BOUND TO TEST LONGER PROGRAMS, OR LOWERING IT
C      TO CONSERVE STORAGE.)                               00001300
C
C      THE OPERATIONS OF THE STRAIGHT-LINE PROGRAM. FORMAT(13,I2,I4).
C      THE I-TH DATA ENTRY IS ENCODED AS I, THE J-TH COMPUTED VALUE AS
C      1ØØ + J AND THE K-TH CONSTANT AS -K. THE OPERATIONS +, -, *, /,
C      SQR AND UNARY MINUS ARE ENCODED AS 1-6, RESPECTIVELY. SQR AND
C      UNARY MINUS REQUIRE Ø AS THEIR SECOND OPERAND.     00001350
C
C      THE NUMBER OF OUTPUTS OF THE STRAIGHT-LINE PROGRAM.
C      GE.1 AND LE.2Ø. FORMAT(I2)                         00001400
C
C      THE INSTRUCTIONS AT WHICH THE OUTPUTS ARE COMPUTED. FORMAT(I3).
C
C      THE NUMBER OF CONSTANTS. LE.2Ø. FORMAT(I2).       00001450
C
C      THE CONSTANTS. FORMAT(G2Ø.16)                     00001500
C
C      THE NUMBER OF ENTRIES IN A SET OF DATA. LE.3Ø. FORMAT(I2).
C
C ----- 00001550
C
C THIS MAIN PROGRAM PERFORMS INPUT AND OUTPUT DUTIES.   00001600
C
C THE SUBPROGRAMS ARE:                                    00001650
C
C      MAXIM - A 'DIRECT SEARCH' NUMERICAL MAXIMIZER CALLED BY THE
C      MAIN PROGRAM.                                       00001700
C
C      GRAM - A GRAM-SCHMIDT ROUTINE USED BY MAXIM.       00001750
C
C      F - FUNCTION CALLED BY MAXIM WHICH EVALUATES THE PENALIZED
C      ERROR-COMPARING VALUE.                              00001800
C
C      ROUND - ROUTINE CALLED BY THE MAIN PROGRAM AND F TO EVALUATE
C      SENSITIVITY TO ERRORS.                              00001850
C
C      GETRHO - ROUTINE CALLED BY ROUND IF THE USER OPTS TO TEST RHO
C      INSTEAD OF OMEGA.                                   00001900
C
C      OMEGA - ROUTINE CALLED BY ROUND TO EVALUATE OMEGA.  00001950
C
C      GENEIG, REDUC, TRED1, TRIDIB - ROUTINES TO COMPUTE THE
C      EIGENVALUES FOR OMEGA. THE USER MAY NEED TO USE DIFFERENT
C      VERSIONS OF THE EISPACK ROUTINES.                   00002000
C
C      POSITV - USER-SUPPLIED ROUTINE TO EVALUATE CONSTRAINTS.
C
C ----- 00002050

```

Contents of program comparing rounding errors in two algorithms:
 MAIN2(354), F(64), ROUND(111), GETRHO(36), OMEGA(40)

C		00000050
C	THE USER SUPPLIES:	00000100
C		00000150
C	THE OUTPUT FROM THE MINICOMPILER FOR EACH OF THE TWO	00000200
C	ALGORITHMS BEING TESTED.	00000250
C		00000300
C	THE ENTRIES OF THE INITIAL SET OF DATA. ONE ENTRY PER CARD, EACH	00000350
C	ENTRY WRITTEN WITH A DECIMAL POINT AND CONTAINED IN THE FIRST	00000400
C	TWENTY COLUMNS.	00000450
C		00000500
C	THE CHOICE OF OMEGA (+) OR RHO (-). FORMAT(I2).	00000550
C	CODE : COMPARE	00000600
C	Ø1 : (ALGORITHM 1) / (ALGORITHM 2)	00000650
C	Ø2 : (ALGORITHM 2) / (ALGORITHM 1)	00000700
C		00000750
C	THE STOPPING VALUE FOR THE MAXIMIZER. WRITTEN WITH A DECIMAL	00000800
C	POINT AND CONTAINED IN THE FIRST TWENTY COLUMNS OF ITS CARD.	00000850
C		00000900
C	THE SOFTWARE RETURNS:	00000950
C		00001000
C	AN ANNOTATED LISTING OF THE USER-SUPPLIED INFORMATION, PLUS THE	00001050
C	ERROR-COMPARING VALUE, THE CONSTRAINT VALUES (IF ANY) AND	00001100
C	THE OUTPUT COMPUTED AT THE INITIAL SET OF DATA.	00001150
C		00001200
C	A LIST OF SELECTED VALUES FOUND BY THE MAXIMIZER.	00001250
C		00001300
C	THE FINAL SET OF DATA.	00001350
C		00001400
C	OTHER INFORMATION IS RETURNED IF EXCEPTIONS ARISE.	00001450

Contents required by both the above rounding error packages and the data for
 the minicompiler: LISTOP(38), MAXIM(109), GRAM(38), GENEIG(44),
 REDUC(119), TRED1(118), TRIDIB(272), MDATA^D(279)

C		00001500
C	-----	00001550
C		00001600
C	THE USER CAN AVOID THE MINICOMPILER BY SUPPLYING:	00001650
C		00001700
C	THE NUMBER OF OPERATIONS IN THE PROGRAM BEING TESTED. GE.1	00001750
C	AND LE.200. FORMAT(I3). (INSTRUCTIONS ARE PROVIDED BELOW FOR	00001800
C	RAISING THE UPPER BOUND TO TEST LONGER PROGRAMS, OR LOWERING IT	00001850
C	TO CONSERVE STORAGE.)	00001900
C		00001950
C	THE OPERATIONS OF THE STRAIGHT-LINE PROGRAM. FORMAT(I3,I2,I4).	00002000
C	THE I-TH DATA ENTRY IS ENCODED AS I, THE J-TH COMPUTED VALUE AS	00002050
C	100 + J AND THE K-TH CONSTANT AS -K. THE OPERATIONS +, -, *, /,	00002100
C	SQRT AND UNARY MINUS ARE ENCODED AS 1-6, RESPECTIVELY. SQRT AND	00002150
C	UNARY MINUS REQUIRE Ø AS THEIR SECOND OPERAND.	00002200
C		00002250
C	THE NUMBER OF OUTPUTS OF THE STRAIGHT-LINE PROGRAM.	00002300
C	GE.1 AND LE.20. FORMAT(I2)	00002350
C		00002400
C	THE INSTRUCTIONS AT WHICH THE OUTPUTS ARE COMPUTED. FORMAT(I3).	00002450
C		00002500
C	THE NUMBER OF CONSTANTS. LE.20. FORMAT(I2).	00002550
C		00002600
C	THE CONSTANTS. FORMAT(G20.16)	00002650
C		00002700
C	THE NUMBER OF ENTRIES IN A SET OF DATA. LE.30. FORMAT(I2).	00002750
C		00002800
C	-----	00002850
C		00002900
C	THIS MAIN PROGRAM PERFORMS INPUT AND OUTPUT DUTIES.	00002950
C		00003000
C	THE SUBPROGRAMS ARE:	00003050
C		00003100
C	MAXIM - A 'DIRECT SEARCH' NUMERICAL MAXIMIZER CALLED BY THE	00003150
C	MAIN PROGRAM.	00003200
C		00003250
C	GRAM - A GRAM-SCHMIDT ROUTINE USED BY MAXIM.	00003300
C		00003350
C	F - FUNCTION CALLED BY MAXIM WHICH EVALUATES THE PENALIZED	00003400

COLLECTED ALGORITHMS (cont.)

532-P 7- 0

C	ERROR-COMPARING VALUE.	00003450
C		00003500
C	ROUND - ROUTINE CALLED BY THE MAIN PROGRAM AND F TO EVALUATE	00003550
C	SENSITIVITY TO ERRORS.	00003600
C		00003650
C	GETRHO - ROUTINE CALLED BY ROUND IF THE USER OPTS TO TEST RHO	00003700
C	INSTEAD OF OMEGA.	00003750
C		00003800
C	OMEGA - ROUTINE CALLED BY ROUND TO EVALUATE OMEGA.	00003850
C		00003900
C	GENEIG, REDUC, TRED1, TRIDIB - ROUTINES TO COMPUTE THE	00003950
C	EIGENVALUES FOR OMEGA. THE USER MAY NEED TO USE DIFFERENT	00004000
C	VERSIONS OF THE EISPACK ROUTINES.	00004050
C		00004100
C	POSITV - USER-SUPPLIED ROUTINE TO EVALUATE CONSTRAINTS.	00004150
C		00004200

ALGORITHM 533

NSPIV, A Fortran Subroutine for Sparse Gaussian Elimination With Partial Pivoting [F4]

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Key Words and Phrases: sparse Gaussian elimination, sparse linear systems, linear equations, partial pivoting algorithms
CR Categories: 5.14
Language: Fortran

DESCRIPTION

1. Introduction

NSPIV is a Fortran subroutine which solves a sparse system of linear equations

$$Ax = b$$

by sparse Gaussian elimination with partial pivoting. More precisely, it performs Gaussian elimination with column interchanges on the nonsingular $N \times N$ matrix A to effectively obtain a factorization of the form

$$AQ = LU,$$

where L is lower triangular, U is unit upper triangular, and Q is a permutation matrix corresponding to the column interchanges. To conserve storage, only the factor U is retained, so during elimination, operations are performed on the right-hand side to obtain the solution y of the system.

$$Ly = b.$$

Once U has been obtained, x is computed by solving the upper triangular system

$$UQ^T x = y.$$

This algorithm discusses the usage of NSPIV and gives a few test results. The method used in NSPIV is described in [7], where it is called "run insertion"; [7] also includes extensive test results which show NSPIV to be more efficient than other currently available software for sparse Gaussian elimination with pivoting.

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2. Usage

For a description of the calling sequence and matrix storage schemes used, see the listing presented here.

The actual numerical computations are performed in an internal subroutine NSPIV1, which is written to perform all computations in single precision. Conversion to double precision may be accomplished simply by changing REAL declarations to DOUBLE PRECISION declarations in both NSPIV and NSPIV1, and by changing the calls to ABS into calls to DABS.

It is well known that the initial ordering of the rows of A greatly affects the overall performance of a subroutine like NSPIV (cf. [1, 7]). To provide initial row and column orderings to NSPIV, the user must set the arrays R , C , and IC so that $R(I)$ is the number of the I th row in the row ordering, $C(I)$ is the number of the I th column in the column ordering, and $IC(C(I)) = I$ for all I . (Notice that no modifications are required to the arrays, IA , JA , and A .) It is always assumed by NSPIV that the user has provided initial row and column orderings; if the orderings are unchanged from those of A , then the user should set $R(I) = C(I) = IC(I) = I$.

Often, sparse linear systems arise with a natural band structure in which all of the nonzeros of A are clustered near the main diagonal. On such problems, with a poor initial ordering, NSPIV may perform worse than some widely available subroutines for band Gaussian elimination with partial pivoting. However, for a sparse matrix A with a given row ordering, NSPIV will often require less storage and not too much more computation time than a band subroutine. Moreover, it is usually possible to find easily computed row orderings for which NSPIV is substantially more efficient than a band subroutine would be with a good band-reducing ordering.

To illustrate these remarks, we consider a sample problem which could arise in the numerical solution of partial differential equations. Here A is a 10×10 block tridiagonal matrix

$$A = \begin{pmatrix} C & D & & & & & & & & \\ & B & C & D & & & & & & \\ & & B & \cdot & \cdot & \cdot & & & & \\ & & & \cdot & \cdot & \cdot & \cdot & & & \\ & & & & \cdot & \cdot & \cdot & & & \\ \cdot & & & & & & & & & D \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & B & C \end{pmatrix}$$

with the 10×10 blocks given by

$$C = \begin{pmatrix} 3. & -0.5 & & & & & & & & \\ -1.5 & 3. & -0.5 & & & & & & & \\ & -1.5 & \cdot & \cdot & \cdot & & & & & \\ & & \cdot & \cdot & \cdot & \cdot & & & & \\ \cdot & & & & & & & & & -0.5 \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & -1.5 & 3. \end{pmatrix}, \quad B = D = \begin{pmatrix} -1. & & & & & & & & & \\ & -1. & & & & & & & & \\ & & \cdot & & & & & & & \\ \cdot & & & \cdot & & & & & & \\ & & & & \cdot & & & & & \\ & & & & & \cdot & & & & \\ & & & & & & \cdot & & & \\ & & & & & & & \cdot & & \\ & & & & & & & & \cdot & \\ & & & & & & & & & -1. \end{pmatrix}$$

We solved this problem using NSPIV with three different initial row orderings: the natural ordering (NO) in which the problem was presented, the so-called alternate diagonal ordering (ADO) (cf. Price and Coats [6]), and an ordering in which the sparsest rows are ordered first (SO). We also solved the problem using the subroutine LEQT1B from the IMSL library ([4]) using the natural ordering to produce the smallest possible bandwidth for this problem (cf. George [3]).

Table I

Subroutine	Ordering	Solution Time*	Total Storage
NSPIV	SO	.799	4785
NSPIV	ADO	.146	2453
NSPIV	NO	.283	3181
LEQT1B	NO	.244	3300

* Times are in seconds. Programs were run on a CDC-6600 computer using the MNF Fortran compiler.

Our results are summarized in Table I. With the SO ordering, NSPIV required much more storage and time than LEQT1B did with the natural ordering. With identical orderings, LEQT1B was slightly faster than NSPIV, but NSPIV required less storage and also saved both the input matrix A and the right-hand side b in their original forms. (LEQT1B does save both triangular factors, however.) Finally, with the ADO ordering, NSPIV used substantially less storage and computation time than LEQT1B. It should be noted in passing that the SO ordering apparently works quite well on problems without the natural band structure present in this example (cf. [1, 7]), even though it did not do well here. Also, experiments indicate that more complex orderings such as the minimum degree or modified Markowitz orderings (cf. Sherman [8]) may be used to good advantage with NSPIV.

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ALGORITHM

NAME(n): indicates a Fortran module with n records
NAME^T: indicates "NAME" is included for testing purposes
NAME^D: indicates "NAME" contains test data
Contents: NSPIV(108), NSPIV1(275), MAIN^T(62), PREORD^T(35),
RESCHK^T(28), PIVCHK^T(35), GENPRB^T(37), DATA^D(960)

```

SUBROUTINE NSPIV (N, IA, JA, A, B, MAX, R, C, IC, X, ITEMP, RTEMP, IERR)
C
C
C NSPIV CALLS NSPIV1 WHICH USES SPARSE GAUSSIAN ELIMINATION WITH
C COLUMN INTERCHANGES TO SOLVE THE LINEAR SYSTEM A X = B. THE
C ELIMINATION PHASE PERFORMS ROW OPERATIONS ON A AND B TO OBTAIN
C A UNIT UPPER TRIANGULAR MATRIX U AND A VECTOR Y. THE SOLUTION
C PHASE SOLVES U X = Y.
C
C
C INPUT ARGUMENTS---
C
C N      INTEGER NUMBER OF EQUATIONS AND UNKNOWNNS
C
C IA     INTEGER ARRAY OF N+1 ENTRIES CONTAINING ROW POINTERS TO A
C        (SEE MATRIX STORAGE DESCRIPTION BELOW)
C

```

C	JA	INTEGER ARRAY WITH ONE ENTRY PER NONZERO IN A, CONTAINING	18
C		COLUMN NUMBERS OF THE NONZEROS OF A. (SEE MATRIX STORAGE	19
C		DESCRIPTION BELOW)	20
C			21
C	A	REAL ARRAY WITH ONE ENTRY PER NONZERO IN A, CONTAINING THE	22
C		ACTUAL NONZEROS. (SEE MATRIX STORAGE DESCRIPTION BELOW)	23
C			24
C	B	REAL ARRAY OF N ENTRIES CONTAINING RIGHT HAND SIDE DATA	25
C			26
C	MAX	INTEGER NUMBER SPECIFYING MAXIMUM NUMBER OF OFF-DIAGONAL	27
C		NONZERO ENTRIES OF U WHICH MAY BE STORED	28
C			29
C	R	INTEGER ARRAY OF N ENTRIES SPECIFYING THE ORDER OF THE	30
C		ROWS OF A (I.E., THE ELIMINATION ORDER FOR THE EQUATIONS)	31
C			32
C	C	INTEGER ARRAY OF N ENTRIES SPECIFYING THE ORDER OF THE	33
C		COLUMNS OF A. C IS ALSO AN OUTPUT ARGUMENT	34
C			35
C	IC	INTEGER ARRAY OF N ENTRIES WHICH IS THE INVERSE OF C	36
C		(I.E., IC(C(I)) = I). IC IS ALSO AN OUTPUT ARGUMENT	37
C			38
C	ITEMP	INTEGER ARRAY OF 2*N + MAX + 2 ENTRIES, FOR INTERNAL USE	39
C			40
C	RTEMP	REAL ARRAY OF N + MAX ENTRIES FOR INTERNAL USE	41
C			42
C			43
C		OUTPUT ARGUMENTS---	44
C			45
C	C	INTEGER ARRAY OF N ENTRIES SPECIFYING THE ORDER OF THE	46
C		COLUMNS OF U. C IS ALSO AN INPUT ARGUMENT	47
C			48
C	IC	INTEGER ARRAY OF N ENTRIES WHICH IS THE INVERSE OF C	49
C		(I.E., IC(C(I)) = I). IC IS ALSO AN INPUT ARGUMENT	50
C			51
C	X	REAL ARRAY OF N ENTRIES CONTAINING THE SOLUTION VECTOR	52
C			53
C	IERR	INTEGER NUMBER WHICH INDICATES ERROR CONDITIONS OR	54
C		THE ACTUAL NUMBER OF OFF-DIAGONAL ENTRIES IN U (FOR	55
C		SUCCESSFUL COMPLETION)	56
C			57
C		IERR VALUES ARE---	58
C			59
C	Ø LT IERR	SUCCESSFUL COMPLETION. U HAS IERR	60
C		OFF-DIAGONAL NONZERO ENTRIES	61
C			62
C	IERR = Ø	ERROR. N = Ø	63
C			64
C	-N LE IERR LT Ø	ERROR. ROW NUMBER IABS(IERR) OF A IS	65
C		IS NULL	66
C			67
C	-2*N LE IERR LT -N	ERROR. ROW NUMBER IABS(IERR+N) HAS A	68
C		DUPLICATE ENTRY	69
C			70
C	-3*N LE IERR LT -2*N	ERROR. ROW NUMBER IABS(IERR+2*N)	71
C		HAS A ZERO PIVOT	72
C			73
C	-4*N LE IERR LT -3*N	ERROR. ROW NUMBER IABS(IERR+3*N)	74
C		EXCEEDS STORAGE	75
C			76
C			77
C		STORAGE OF SPARSE MATRICES---	78
C			79
C		THE SPARSE MATRIX A IS STORED USING THREE ARRAYS IA, JA, AND A.	80
C		THE ARRAY A CONTAINS THE NONZEROS OF THE MATRIX ROW-BY-ROW, NOT	81
C		NECESSARILY IN ORDER OF INCREASING COLUMN NUMBER. THE ARRAY JA	82
C		CONTAINS THE COLUMN NUMBERS CORRESPONDING TO THE NONZEROS STORED	83
C		IN THE ARRAY A (I.E., IF THE NONZERO STORED IN A(K) IS IN	84
C		COLUMN J, THEN JA(K) = J). THE ARRAY IA CONTAINS POINTERS TO THE	85
C		ROWS OF NONZEROS/COLUMN INDICES IN THE ARRAY A/JA (I.E.,	86
C		A(IA(I))/JA(IA(I)) IS THE FIRST ENTRY FOR ROW I IN THE ARRAY A/JA).	87
C		IA(N+1) IS SET SO THAT IA(N+1) - IA(1) = THE NUMBER OF NONZEROS IN A	88
C			89
C			90
C		REAL A(1),B(1),X(1),RTEMP(1)	91
C		INTEGER IA(1),JA(1),R(1),C(1),IC(1),ITEMP(1)	92
C		INTEGER IU,JU,U,Y,P	93

```

C 94
C SET INDICES TO DIVIDE TEMPORARY STORAGE FOR NSPIV1 95
C 96
C   Y = 1 97
C   U = Y + N 98
C   P = 1 99
C   IU = P + N + 1 100
C   JU = IU + N + 1 101
C 102
C CALL NSPIV1 TO PERFORM COMPUTATIONS 103
C 104
C   CALL NSPIV1 (N, IA, JA, A, B, MAX, R, C, IC, X, RTEMP(Y), ITEMP(P),
C   ITEMP(IU), ITEMP(JU), RTEMP(U), IERR) 105
C RETURN 106
C END 107
C SUBROUTINE NSPIV1 (N, IA, JA, A, B, MAX, R, C, IC, X, Y, P, IU, JU, U, IERR) 108
C 109
C 110
C 111
C NSPIV1 USES SPARSE GAUSSIAN ELIMINATION WITH 112
C COLUMN INTERCHANGES TO SOLVE THE LINEAR SYSTEM  $A X = B$ . THE 113
C ELIMINATION PHASE PERFORMS ROW OPERATIONS ON A AND B TO OBTAIN 114
C A UNIT UPPER TRIANGULAR MATRIX U AND A VECTOR Y. THE SOLUTION 115
C PHASE SOLVES  $U X = Y$ . 116
C 117
C 118
C SEE NSPIV FOR DESCRIPTIONS OF ALL INPUT AND OUTPUT ARGUMENTS 119
C OTHER THAN THOSE DESCRIBED BELOW 120
C 121
C INPUT ARGUMENTS (USED INTERNALLY ONLY)--- 122
C 123
C Y REAL ARRAY OF N ENTRIES USED TO COMPUTE THE UPDATED 124
C RIGHT HAND SIDE 125
C 126
C P INTEGER ARRAY OF N+1 ENTRIES USED FOR A LINKED LIST. 127
C P(N+1) IS THE LIST HEADER, AND THE ENTRY FOLLOWING 128
C P(K) IS IN P(P(K)). THUS, P(N+1) IS THE FIRST DATA 129
C ITEM, P(P(N+1)) IS THE SECOND, ETC. A POINTER OF 130
C N+1 MARKS THE END OF THE LIST 131
C 132
C IU INTEGER ARRAY OF N+1 ENTRIES USED FOR ROW POINTERS TO U 133
C (SEE MATRIX STORAGE DESCRIPTION BELOW) 134
C 135
C JU INTEGER ARRAY OF MAX ENTRIES USED FOR COLUMN NUMBERS OF 136
C THE NONZEROS IN THE STRICT UPPER TRIANGLE OF U. (SEE 137
C MATRIX STORAGE DESCRIPTION BELOW) 138
C 139
C U REAL ARRAY OF MAX ENTRIES USED FOR THE ACTUAL NONZEROS IN 140
C THE STRICT UPPER TRIANGLE OF U. (SEE MATRIX STORAGE 141
C DESCRIPTION BELOW) 142
C 143
C 144
C STORAGE OF SPARSE MATRICES--- 145
C 146
C THE SPARSE MATRIX A IS STORED USING THREE ARRAYS IA, JA, AND A. 147
C THE ARRAY A CONTAINS THE NONZEROS OF THE MATRIX ROW-BY-ROW, NOT 148
C NECESSARILY IN ORDER OF INCREASING COLUMN NUMBER. THE ARRAY JA 149
C CONTAINS THE COLUMN NUMBERS CORRESPONDING TO THE NONZEROS STORED 150
C IN THE ARRAY A (I.E., IF THE NONZERO STORED IN A(K) IS IN 151
C COLUMN J, THEN JA(K) = J). THE ARRAY IA CONTAINS POINTERS TO THE 152
C ROWS OF NONZEROS/COLUMN INDICES IN THE ARRAY A/JA (I.E., 153
C  $A(IA(I))/JA(IA(I))$  IS THE FIRST ENTRY FOR ROW I IN THE ARRAY A/JA). 154
C IA(N+1) IS SET SO THAT  $IA(N+1) - IA(1) =$  THE NUMBER OF NONZEROS IN 155
C A. IU, JU, AND U ARE USED IN A SIMILAR WAY TO STORE THE STRICT UPPER 156
C TRIANGLE OF U, EXCEPT THAT JU ACTUALLY CONTAINS C(J) INSTEAD OF J 157
C 158
C 159
C REAL A(1),B(1),U(1),X(1),Y(1) 160
C REAL DK,LKI,ONE,XPV,XPVMAX,YK,ZERO 161
C INTEGER C(1),IA(1),IC(1),IU(1),JA(1),JU(1),P(1),R(1) 162
C INTEGER CK,PK,PPK,PV,V,VI,VJ,VK 163
C 164
C 165
C IF (N .EQ. 0) GO TO 1001 166
C 167
C ONE = 1.0 168
C ZERO = 0.0 169
C 170

```

```

C INITIALIZE WORK STORAGE AND POINTERS TO JU          171
C                                                    172
C      DO 10 J=1,N                                  173
C          X(J) = ZERO                                174
10    CONTINUE                                       175
C          IU(1) = 1                                  176
C          JUPTR = 0                                  177
C                                                    178
C PERFORM SYMBOLIC AND NUMERIC FACTORIZATION ROW BY ROW 179
C VK (VI,VJ) IS THE GRAPH VERTEX FOR ROW K (I,J) OF U 180
C                                                    181
C      DO 170 K=1,N                                  182
C                                                    183
C INITIALIZE LINKED LIST AND FREE STORAGE FOR THIS ROW 184
C THE R(K)-TH ROW OF A BECOMES THE K-TH ROW OF U.    185
C                                                    186
C          P(N+1) = N+1                               187
C          VK = R(K)                                  188
C                                                    189
C SET UP ADJACENCY LIST FOR VK, ORDERED IN           190
C CURRENT COLUMN ORDER OF U. THE LOOP INDEX         191
C GOES DOWNWARD TO EXPLOIT ANY COLUMNS             192
C FROM A IN CORRECT RELATIVE ORDER                 193
C                                                    194
C          JMIN = IA(VK)                              195
C          JMAX = IA(VK+1) - 1                        196
C          IF (JMIN .GT. JMAX) GO TO 1002            197
C          J = JMAX                                    198
20    JAJ = JA(J)                                     199
C          VJ = IC(JAJ)                               200
C                                                    201
C STORE A(K,J) IN WORK VECTOR                       202
C                                                    203
C          X(VJ) = A(J)                               204
C THIS CODE INSERTS VJ INTO ADJACENCY LIST OF VK    205
C          PPK = N+1                                  206
30    PPK = PPK                                       207
C          PPK = P(PK)                                208
C          IF (PPK - VJ) 30,1003,40                 209
40    P(VJ) = PPK                                     210
C          P(PK) = VJ                                 211
C          J = J - 1                                  212
C          IF (J .GE. JMIN) GO TO 20                 213
C                                                    214
C THE FOLLOWING CODE COMPUTES THE K-TH ROW OF U     215
C                                                    216
C          VI = N+1                                    217
C          YK = B(VK)                                  218
50    VI = P(VI)                                       219
C          IF (VI .GE. K) GO TO 110                 220
C                                                    221
C VI LT VK -- PROCESS THE L(K,I) ELEMENT AND MERGE THE 222
C ADJACENCY OF VI WITH THE ORDERED ADJACENCY OF VK 223
C                                                    224
C          LKI = - X(VI)                              225
C          X(VI) = ZERO                               226
C                                                    227
C ADJUST RIGHT HAND SIDE TO REFLECT ELIMINATION    228
C                                                    229
C          YK = YK + LKI * Y(VI)                    230
C          PPK = VI                                    231
C          JMIN = IU(VI)                              232
C          JMAX = IU(VI+1) - 1                        233
C          IF (JMIN .GT. JMAX) GO TO 50             234
C          DO 100 J=JMIN,JMAX                         235
C              JUJ = JU(J)                            236
C              VJ = IC(JUJ)                           237
C                                                    238
C IF VJ IS ALREADY IN THE ADJACENCY OF VK,         239
C SKIP THE INSERTION                                240
C                                                    241
C          IF (X(VJ) .NE. ZERO) GO TO 90           242
C                                                    243
C INSERT VJ IN ADJACENCY LIST OF VK.              244
C RESET PPK TO VI IF WE HAVE PASSED THE CORRECT    245
C INSERTION SPOT. (THIS HAPPENS WHEN THE ADJACENCY OF 246

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```

C VI IS NOT IN CURRENT COLUMN ORDER DUE TO PIVOTING.) 247
C 248
      IF (VJ - PPK) 60,90,70 249
60   PPK = VI 250
70   PK = PPK 251
      PPK = P(PK) 252
      IF (PPK - VJ) 70,90,80 253
80   P(VJ) = PPK 254
      P(PK) = VJ 255
      PPK = VJ 256
C 257
C COMPUTE L(K,J) = L(K,J) - L(K,I)*U(I,J) FOR L(K,I) NONZERO 258
C COMPUTE U*(K,J) = U*(K,J) - L(K,I)*U(I,J) FOR U(K,J) NONZERO 259
C (U*(K,J) = U(K,J)*D(K,K)) 260
C 261
90   X(VJ) = X(VJ) + LKI * U(J) 262
100  CONTINUE 263
      GO TO 50 264
C 265
C PIVOT--INTERCHANGE LARGEST ENTRY OF K-TH ROW OF U WITH 266
C THE DIAGONAL ENTRY. 267
C 268
C FIND LARGEST ENTRY, COUNTING OFF-DIAGONAL NONZEROES 269
C 270
110  IF (VI .GT. N) GO TO 1004 271
      XPVMAX = ABS(X(VI)) 272
      MAXC = VI 273
      NZCNT = 0 274
      PV = VI 275
120  V = PV 276
      PV = P(PV) 277
      IF (PV .GT. N) GO TO 130 278
      NZCNT = NZCNT + 1 279
      XPV = ABS(X(PV)) 280
      IF (XPV .LE. XPVMAX) GO TO 120 281
      XPVMAX = XPV 282
      MAXC = PV 283
      MAXCL = V 284
      GO TO 120 285
130  IF (XPVMAX .EQ. ZERO) GO TO 1004 286
C 287
C IF VI = K, THEN THERE IS AN ENTRY FOR DIAGONAL 288
C WHICH MUST BE DELETED. OTHERWISE, DELETE THE 289
C ENTRY WHICH WILL BECOME THE DIAGONAL ENTRY 290
C 291
      IF (VI .EQ. K) GO TO 140 292
      IF (VI .EQ. MAXC) GO TO 140 293
      P(MAXCL) = P(MAXC) 294
      GO TO 150 295
140  VI = P(VI) 296
C 297
C COMPUTE D(K) = 1/L(K,K) AND PERFORM INTERCHANGE. 298
C 299
150  DK = ONE / X(MAXC) 300
      X(MAXC) = X(K) 301
      I = C(K) 302
      C(K) = C(MAXC) 303
      C(MAXC) = I 304
      CK = C(K) 305
      IC(CK) = K 306
      IC(I) = MAXC 307
      X(K) = ZERO 308
C 309
C UPDATE RIGHT HAND SIDE. 310
C 311
      Y(K) = YK * DK 312
C 313
C COMPUTE VALUE FOR IU(K+1) AND CHECK FOR STORAGE OVERFLOW 314
C 315
      IU(K+1) = IU(K) + NZCNT 316
      IF (IU(K+1) .GT. MAX+1) GO TO 1005 317
C 318
C MOVE COLUMN INDICES FROM LINKED LIST TO JU. 319
C COLUMNS ARE STORED IN CURRENT ORDER WITH ORIGINAL 320
C COLUMN NUMBER (C(J)) STORED FOR CURRENT COLUMN J 321
C 322

```

```

        IF (VI .GT. N) GO TO 170
        J = VI
160    JUPTR = JUPTR + 1
        JU(JUPTR) = C(J)
        U(JUPTR) = X(J) * DK
        X(J) = 'ZERO
        J = P(J)
        IF (J .LE. N) GO TO 160
170    CONTINUE
C
C BACKSOLVE U X = Y, AND REORDER X TO CORRESPOND WITH A
C
    K = N
    DO 200 I=1,N
        YK = Y(K)
        JMIN = IU(K)
        JMAX = IU(K+1) - 1
        IF (JMIN .GT. JMAX) GO TO 190
        DO 180 J=JMIN,JMAX
            JUJ = JU(J)
            JUJ = IC(JUJ)
            YK = YK - U(J) * Y(JUJ)
180    CONTINUE
190    Y(K) = YK
        CK = C(K)
        X(CK) = YK
        K = K-1
200    CONTINUE
C
C RETURN WITH IERR = NUMBER OF OFF-DIAGONAL NONZEROES IN U
C
    IERR = IU(N+1) - IU(1)
    RETURN
C
C ERROR RETURNS
C
C N = 0
C
1001 IERR = 0
    RETURN
C
C ROW K OF A IS NULL
C
1002 IERR = -K
    RETURN
C
C ROW K OF A HAS A DUPLICATE ENTRY
C
1003 IERR = -(N+K)
    RETURN
C
C ZERO PIVOT IN ROW K
C
1004 IERR = -(2*N+K)
    RETURN
C
C STORAGE FOR U EXCEEDED ON ROW K
C
1005 IERR = -(3*N+K)
    RETURN
    END

```

APPENDIX A

The code in this appendix illustrates the use of NSPIV. The linear system is of the form

$$Ax = b$$

where A is a 10×10 block tridiagonal matrix with 10×10 blocks. Specifically, This example is chosen for its simplicity; it does not exercise the algorithm, since A is a strictly diagonally dominant matrix.

$$A = \begin{pmatrix} C & D & & & \\ B & C & D & \bigcirc & \\ & B & \cdot & \cdot & \cdot \\ & & \cdot & \cdot & \cdot \\ \bigcirc & & \cdot & \cdot & D \\ & & & B & C \end{pmatrix}$$

with

$$B = \begin{pmatrix} -1 & & & & \\ & -1 & & \bigcirc & \\ & & \cdot & \cdot & \\ \bigcirc & & & \cdot & \\ & & & & -1 \end{pmatrix}, \quad C = \begin{pmatrix} & & & & \\ -1 & 4 & & & \\ & -1 & \cdot & & \bigcirc \\ & & \cdot & \cdot & \cdot \\ \bigcirc & & \cdot & \cdot & \\ & & & -1 & 4 \end{pmatrix}$$

$$D = \begin{pmatrix} -1.5 & & & & \\ & -1.5 & & \bigcirc & \\ & & \cdot & \cdot & \\ \bigcirc & & & \cdot & \\ & & & & -1.5 \end{pmatrix}$$

```

C      PROGRAM PIVCHK(OUTPUT,TAPE6=OUTPUT)
C
C      THIS PROGRAM ILLUSTRATES THE USE OF NSPIV BY SOLVING THE
C      SYSTEM OF LINEAR EQUATIONS
C
C      A X = B
C
C      WITH A AN NG X NG BLOCK TRIDIAGONAL MATRIX, WITH NG X NG BLOCKS.
C      THE DIAGONAL BLOCKS OF A ARE LOWER BI-DIAGONAL (ENTRIES ARE 4.0
C      ON THE DIAGONAL, -1.0 ON THE SUBDIAGONAL), AND THE OFF-DIAGONAL
C      BLOCKS OF A ARE DIAGONAL (ENTRIES ARE -1.0 IN THE LOWER TRIANGLE,
C      -1.5 IN THE UPPER TRIANGLE.) X IS CHOSEN TO BE A VECTOR
C      OF ALL ONES, AND B IS COMPUTED ACCORDINGLY.
C
C      INTEGER IA(101),JA(400),R(100),C(100),IC(100),ITEMP(597)
C      REAL A(400),B(100),X(100),RTEMP(495)
C      DATA MAX/395/,NG/10/,N/100/
C
C      SET UP PROBLEM
C
C      K = 1
C      IA(1) = 1
C      IAPTR = 1
C      DO 5 I=1,NG
C        DO 5 J=1,NG
C          BK = 0.
C          IF (I .EQ. 1) GO TO 1
C          JA(IAPTR) = K - NG
C          A(IAPTR) = -1.
C          BK = BK - 1.
C          IAPTR = IAPTR + 1
C          IF (J .EQ. 1) GO TO 2
C          JA(IAPTR) = K - 1
C          A(IAPTR) = -1.
C          BK = BK - 1.
C
1

```

```

      IAPTR = IAPTR + 1
2     JA(IAPTR) = K
      A(IAPTR) = 4.
      BK = BK + 4.
      IAPTR = IAPTR + 1
      IF (I .EQ. NG) GO TO 4
      JA(IAPTR) = K + NG
      A(IAPTR) = -1.5
      BK = BK - 1.5
      IAPTR = IAPTR + 1
4     B(K) = BK
      K = K + 1
      IA(K) = IAPTR
5     CONTINUE
C
C CALL PREORD TO ORDER ROWS OF A BY INCREASING NUMBERS OF NONZEROES
C
      CALL PREORD(N,IA,R,C,IC)
C
C CALL NSPIV TO SOLVE SYSTEM
C
      CALL NSPIV(N,IA,JA,A,B,MAX,R,C,IC,X,ITEMP,RTEMP,IERR)
101  WRITE (6,101) IERR
101  FORMAT (8H IERR = ,I10)
C
C CALL RESCHK TO COMPUTE MAX-NORM AND 2-NORM OF RESIDUAL
C
      CALL RESCHK(N,IA,JA,A,B,X)
C
      STOP
      END

      SUBROUTINE PREORD(N,IA,R,C,IC)
C
C PREORD ORDERS THE ROWS OF A BY INCREASING NUMBER OF NONZEROES.
C THE ROW PERMUTATION IS RETURNED IN R. C IS SET TO THE IDENTITY.
C
      INTEGER IA(1),R(1),C(1),IC(1)
C
      DO 1 I=1,N
          R(I) = I
          C(I) = I
          IC(I) = I
1     CONTINUE
      DO 5 I = 1,N
          C(I) = 0
5     DO 10 K = 1,N
          KDEG = IA(K+1) - IA(K)
          IF (KDEG .EQ. 0) KDEG = KDEG + 1
          IC(K) = C(KDEG)
          C(KDEG) = K
10    CONTINUE
          I = 0
      DO 30 J = 1,N
          IF (C(J) .EQ. 0) GO TO 30
          K = C(J)
20    I = I + 1
          R(I) = K
          K = IC(K)
          IF (K .GT. 0) GO TO 20
30    CONTINUE
      DO 40 I = 1,N
          C(I) = I
          IC(I) = I
40    CONTINUE
      RETURN
      END
      SUBROUTINE RESCHK(N,IA,JA,A,B,X)
C
C RESCHK COMPUTES THE MAX-NORM AND 2-NORM OF THE RESIDUAL.
C DOUBLE PRECISION IS USED FOR THE COMPUTATION.
C
      INTEGER IA(1),JA(1)
      REAL A(1),B(1),X(1)
      DOUBLE PRECISION RESID,RESIDM,ROWSUM
      RESID = 0.

```



```

RESIDM = 0.
DO 20 I=1,N
  ROWSUM = DBLE(B(I))
  JMIN = IA(I)
  JMAX = IA(I+1) - 1
  DO 10 J=JMIN,JMAX
    JAJ = JA(J)
    ROWSUM = ROWSUM - DBLE(A(J)) * DBLE(X(JAJ))
10  CONTINUE
    IF (DABS(ROWSUM) .GT. RESIDM) RESIDM = DABS(ROWSUM)
    RESID = RESID + ROWSUM**2
20  CONTINUE
RESID = DSQRT(RESID)
WRITE (6,25) RESID
25  FORMAT (22H 2-NORM OF RESIDUAL = ,D14.7)
WRITE (6,30) RESIDM
30  FORMAT (24H MAX NORM OF RESIDUAL = ,D14.7)
RETURN
END
112
113
114
115
116
117
118
119
120
121
122
123
124
125
126
127
128
129
130

C   PROGRAM PIVCHK(INPUT,OUTPUT,TAPE6=OUTPUT,TAPE5=INPUT)
C
C   THIS PROGRAM ILLUSTRATES THE USE OF NSPIV BY SOLVING THE
C   SYSTEM OF LINEAR EQUATIONS
C
C       A X = B
C
C   WHERE A IS A BANDED 192 X 192 MATRIX OF BANDWIDTH 20. X IS
C   CHOSEN TO BE A VECTOR OF ALL ONES, AND B IS COMPUTED ACCORDINGLY.
C
C   INTEGER IA(193),JA(3500),R(192),C(192),IC(192),ITEMP(7500)
C   REAL A(3500),B(192),X(192),RTEMP(7200)
C   DATA MAX/7000/,N/192/
C
C   CALL GENPRB TO SET UP PROBLEM
C
C       CALL GENPRB(N,IA,JA,A,ITEMP,B)
C
C   CALL PREORD TO ORDER ROWS OF A BY INCREASING NUMBERS OF NONZEROES
C
C       CALL PREORD(N,IA,R,C,IC)
C
C   CALL NSPIV TO SOLVE SYSTEM
C
C       CALL NSPIV(N,IA,JA,A,B,MAX,R,C,IC,X,ITEMP,RTEMP,IERR)
C       WRITE (6,101) IERR
101  FORMAT (8H IERR = ,I10)
C
C   CALL RESCHK TO COMPUTE MAX-NORM AND 2-NORM OF RESIDUAL
C
C       CALL RESCHK(N,IA,JA,A,B,X)
C
C   STOP
C   END
C   SUBROUTINE GENPRB(N,IA,JA,A,B,RHS)
C
C   GENPRB SETS UP THE MATRIX AND RIGHT HAND SIDE FROM
C   THE DATA ON CARDS. THE PROBLEM IS ONE FROM THE USGS.
C
C   INTEGER IA(1),JA(1)
C   REAL B(1),A(1),RHS(1)
C
C       KMIN = 1
C       DO 5 I=1,N
C         KMAX = KMIN + 38
C         READ (5,101) (B(K),K=KMIN,KMAX)
101  FORMAT (8G10.3)
C         KMIN = KMAX + 1
5     CONTINUE
C
C       IA(1) = 1
C       IAPTR = 1
C       IPART = 0
C       DO 15 I=1,N
C         RHSI = 0.

```

```

DO 10 J=1, 39
K = I - 20 + J
IF (K .LE. 0) GO TO 10
IF (K .GT. N) GO TO 10
IPARTJ = IPART + J
IF (B(IPARTJ) .EQ. 0.) GO TO 10
A(IPATR) = B(IPARTJ)
JA(IPATR) = K
RHSI = RHSI + B(IPATR)
IPATR = IPATR + 1
10 CONTINUE
IA(I+1) = IPATR
RHS(1) = RHSI
15 IPART = IPART + 39
RETURN
END
0 0 0 0 0 0 0
0 0 0 0 0 0 0
0 0 0 2.67 - .819 - .475 - .409 0
0 0 0 - .475 - .409 - 1.14 - .204 0
0 0 0 0 0 0 0 0
0 0 0 0 0 0 0 0
0 0 1.960E-02 .125 -7.451E-03 6.258E-02 0 0
0 0 -7.451E-03 6.258E-02 -1.233E-02 3.117E-02 0 0
0 0 0 0 0 0 0 0
0 -2.970E-02 -2.556E-02 .333 - .102 -2.970E-02 -2.556E-02 0
0 0 0 0 0 0 0 0
0 -7.142E-02 -1.278E-02 -5.939E-02 -5.120E-02 -7.142E-02 -1.278E-02 0
0 0 0 0 0 0 0 0
0 -7.451E-03 6.258E-02 3.912E-02 .249 -7.451E-03 6.258E-02 0
0 0 0 0 0 0 0 0
-1.233E-02 3.117E-02 -1.490E-02 .125 -1.233E-02 3.117E-02 0 0
0 0 0 0 0 0 0 0
0 -2.970E-02 -2.556E-02 .333 - .102 -2.970E-02 -2.556E-02 0
0 0 0 0 0 0 0 0
0 -7.142E-02 -1.278E-02 -5.939E-02 -5.120E-02 -7.142E-02 -1.278E-02 0
0 0 0 0 0 0 0 0
- .119 1.000 .626 3.99 - .119 1.000 0 0
- .197 .499 - .238 1.99 - .197 .499 0 0
0 0 0 0 0 0 0 0
0 -2.970E-02 -2.556E-02 .333 - .102 -2.970E-02 -2.556E-02 0
0 0 0 0 0 0 0 0
0 -7.142E-02 -1.278E-02 -5.939E-02 -5.120E-02 -7.142E-02 -1.278E-02 0
0 0 0 0 0 0 0 0
- .119 1.000 .626 3.99 - .119 1.000 0 0
- .197 .499 - .238 1.99 - .197 .499 0 0
0 0 0 0 0 0 0 0
0 -2.970E-02 -2.556E-02 .333 - .102 -2.970E-02 -2.556E-02 0
0 0 0 0 0 0 0 0
0 -7.142E-02 -1.278E-02 -5.939E-02 -5.120E-02 -7.142E-02 -1.278E-02 0
0 0 0 0 0 0 0 0
- .119 1.000 .626 3.99 - .119 1.000 0 0
- .197 .499 - .238 1.99 - .197 .499 0 0
0 0 0 0 0 0 0 0
0 -2.970E-02 -2.556E-02 .333 - .102 -2.970E-02 -2.556E-02 0
0 0 0 0 0 0 0 0
0 -7.142E-02 -1.278E-02 -5.939E-02 -5.120E-02 -7.142E-02 -1.278E-02 0
0 0 0 0 0 0 0 0
- .119 1.000 .626 3.99 - .119 1.000 0 0
- .197 .499 - .238 1.99 - .197 .499 0 0

```



```

0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0
0-5.939E-02-5.120E-02 .668 -.205 -5.939E-02-5.120E-02 0
0 0 0 0 0 0 0 0
-.197 0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
.499 -.238 1.99 -.197 .499 0 0 0
0 0 0 0 0 0 0 0
-.238 1.99 1.25 7.98 -.238 1.99 0 0
0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0
0-5.939E-02-5.120E-02 .668 -.205 -5.939E-02-5.120E-02 0
0 0 0 0 0 0 0 0
-.197 0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
.499 -.238 1.99 -.197 .499 0 0
0 0 0 0 0 0 0 0
-.238 1.99 1.25 7.98 -.238 1.99 0 0
0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0
0-5.939E-02-5.120E-02 .668 -.205 -5.939E-02-5.120E-02 0
0 0 0 0 0 0 0 0
-.197 0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
.499 -.238 1.99 -.197 .499 0 0
0 0 0 0 0 0 0 0
-.238 1.99 1.25 7.98 -.238 1.99 0 0
0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0
0-7.142E-02-1.278E-02-2.970E-02-2.556E-02 0 0 0 0
0 0 0 0 0 0 0 0
0-5.939E-02-5.120E-02 .333 -.102 0 0 0 0
0 0 0 0 0 0 0 0
-.197 0-7.142E-02-1.278E-02-2.970E-02-2.556E-02 0 0 0 0
.499 -.119 1.00 0 0 0 0 0
0 0 0 0 0 0 0 0
-.238 1.99 .626 3.99 0 0 0 0
0 0 0 0 0 0 0 0
-.197 .499 -.119 1.00 0 0 0 0
0 0 0 0-2.970E-02-2.556E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0
0 0 0 .333 -.102 -5.939E-02-5.120E-02 0
0 0 0 0 0 0 0 0
0 0 0-2.970E-02-2.556E-02-7.142E-02-1.278E-02 0
0 0 -.119 1.00 -.197 .499 0 0
0 0 0 0 0 0 0 0
0 0 .626 3.99 -.238 1.99 0 0
0 0 0 0 0 0 0 0
0 0 -.119 1.00 -.197 .499 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0
0-5.939E-02-5.120E-02 .668 -.205 -5.939E-02-5.120E-02 0
0 0 0 0 0 0 0 0
-.197 0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
.499 -.238 1.99 -.197 .499 0 0
0 0 0 0 0 0 0 0
-.238 1.99 1.25 7.98 -.238 1.99 0 0
0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0
0-5.939E-02-5.120E-02 .668 -.205 -5.939E-02-5.120E-02 0
0 0 0 0 0 0 0 0
-.197 0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
.499 -.238 1.99 -.197 .499 0 0
0 0 0 0 0 0 0 0
-.238 1.99 1.25 7.98 -.238 1.99 0 0
0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0
0-5.939E-02-5.120E-02 .668 -.205 -5.939E-02-5.120E-02 0
0 0 0 0 0 0 0 0
-.197 0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
.499 -.238 1.99 -.197 .499 0 0
0 0 0 0 0 0 0 0

```



```
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02
-.197 .499 -.238 1.99 -.197 .499 0 0 0
0 0 0 0 0 0 0 0 0
-.238 1.99 1.25 7.98 -.238 1.99 0 0 0
0 0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0 0
0-5.939E-02-5.120E-02 .668 -.205 -5.939E-02-5.120E-02 0
0 0 0 0 0 0 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02
-.197 .499 -.238 1.99 -.197 .499 0 0 0
0 0 0 0 0 0 0 0 0
-.238 1.99 1.25 7.98 -.238 1.99 0 0 0
0 0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0 0
0-5.939E-02-5.120E-02 .668 -.205 -5.939E-02-5.120E-02 0
0 0 0 0 0 0 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02
-.197 .499 -.238 1.99 -.197 .499 0 0 0
0 0 0 0 0 0 0 0 0
-.238 1.99 1.25 7.98 -.238 1.99 0 0 0
0 0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0 0
0-5.939E-02-5.120E-02 .668 -.205 -5.939E-02-5.120E-02 0
0 0 0 0 0 0 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02
-.197 .499 -.238 1.99 -.197 .499 0 0 0
0 0 0 0 0 0 0 0 0
-.238 1.99 1.25 7.98 -.238 1.99 0 0 0
0 0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0 0
0-7.142E-02-1.278E-02-2.970E-02-2.556E-02 0 0 0 0 0
0 0 0 0 0 0 0 0 0
0-5.939E-02-5.120E-02 .333 -.102 0 0 0 0 0
0 0 0 0 0 0 0 0 0
0-7.142E-02-1.278E-02-2.970E-02-2.556E-02 0 0 0 0 0
-.197 .499 -.119 1.00 0 0 0 0 0
0 0 0 0 0 0 0 0 0
-.238 1.99 .626 3.99 0 0 0 0 0
0 0 0 0 0 0 0 0 0
-.197 .499 -.119 1.00 0 0 0 0 0
0 0 0 0-2.970E-02-2.556E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0 0
0 0 0 .333 -.102 -5.939E-02-5.120E-02 0
0 0 0 0 0 0 0 0 0
0-2.970E-02-2.556E-02-7.142E-02-1.278E-02 0
0 0 -.119 1.00 -.197 .499 0 0 0
0 0 0 0 0 0 0 0 0
0 0 .626 3.99 -.238 1.99 0 0 0
0 0 0 0 0 0 0 0 0
0 0 -.119 1.00 -.197 .499 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0 0
0-5.939E-02-5.120E-02 .668 -.205 -5.939E-02-5.120E-02 0
0 0 0 0 0 0 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02
-.197 .499 -.238 1.99 -.197 .499 0 0 0
0 0 0 0 0 0 0 0 0
-.238 1.99 1.25 7.98 -.238 1.99 0 0 0
0 0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0 0
0-5.939E-02-5.120E-02 .668 -.205 -5.939E-02-5.120E-02 0
0 0 0 0 0 0 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02
-.197 .499 -.238 1.99 -.197 .499 0 0 0
0 0 0 0 0 0 0 0 0
-.238 1.99 1.25 7.98 -.238 1.99 0 0 0
0 0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0 0
```



```
-.197 .499 -.238 1.99 -.197 .499 0 0
0 0 0 0 0 0 0 0
-.119 1.00 .626 3.99 -.119 1.00 0 0
0 0 0 0 0 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0
0-2.970E-02-2.556E-02 .333 -.102 -2.970E-02-2.556E-02 0
0 0 0 0 0 0 0 0
0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0
0 0 0 0 0 0 0 0
-.119 1.00 .626 3.99 -.119 1.00 0 0
0 0 0 0 0 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0
0-2.970E-02-2.556E-02 .333 -.102 -2.970E-02-2.556E-02 0
0 0 0 0 0 0 0 0
0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0
0 0 0 0 0 0 0 0
-.119 1.00 .626 3.99 -.119 1.00 0 0
0 0 0 0 0 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0
0-2.970E-02-2.556E-02 .333 -.102 -2.970E-02-2.556E-02 0
0 0 0 0 0 0 0 0
0 0 0 0 0 0 0 0
-.197 .499 -.238 1.99 -.197 .499 0 0
0 0 0 0 0 0 0 0
-.119 1.00 .626 3.99 -.119 1.00 0 0
0 0 0 0 0 0 0 0
0-7.142E-02-1.278E-02-5.939E-02-5.120E-02-7.142E-02-1.278E-02 0
0 0 0 0 0 0 0 0
0-2.970E-02-2.556E-02 .333 -.102 -2.970E-02-2.556E-02 0
0 0 0 0 0 0 0 0
0 0 0 0 0 0 0 0
-1.233E-02 3.117E-02-1.490E-02 .125 -1.233E-02 3.117E-02 0 0
0 0 0 0 0 0 0 0
-7.451E-03 6.258E-02 3.912E-02 .249 -7.451E-03 6.258E-02 0 0
0 0 0 0 0 0 0 0
0 -1.14 -.204 -.475 -.409 0 0 0
0 0 0 0 0 0 0 0
0 -.475 -.409 2.67 -.819 0 0 0
0 0 0 0 0 0 0 0
0 0 0 0 0 0 0 0
-1.233E-02 3.117E-02-7.451E-03 6.258E-02 0 0 0 0
0 0 0 0 0 0 0 0
-7.451E-03 6.258E-02 1.960E-02 .125 0 0 0 0
0 0 0 0 0 0 0 0
0 0 0 0 0 0 0 0
```

ALGORITHM 534

STINT: STiff (differential equations)

INTEgrator [D2]

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 IBM Poughkeepsie Laboratories
 THEODORE A. BICKART
 Syracuse University
 and
 ZDENEK PICEL

Key Words and Phrases: stiff differential equations, stiffly stable methods, composite multistep methods, cyclic methods, numerical integration, ordinary differential equations, initial value problems, multistep formulas, numerical integration program, Fortran code STINT
 CR Categories: 5.16, 5.17
 Language: Fortran

DESCRIPTION

The algorithm given here is a complement to [1] where the theoretical development and test results are described.

REFERENCES

[1] TENDLER, J.M., BICKART, T.A., AND PICEL, Z. A stiffly stable integration process using cyclic composite methods. *ACM Trans. Math. Software* 4, 4 (Dec. 1978), 339-368.

ALGORITHM

```

C*                                *MAIN 10
C*----- THIS IS A MAIN PROGRAM FOR STINT-TYPE SUBROUTINES -----*MAIN 20
C*----- SET HERE FOR N = 4 PROBLEM - -----*MAIN 30
C*                                *MAIN 40
      DOUBLE PRECISION D, EPS, ERO, ERORI, HI, HMAX, HMIN, HNEXT,      MAIN 50
      1                H0, RJ, RW, S, SAVE, T, TF, TI, TOUT,          MAIN 60
      2                TOUTP, TS, YDOT, YI, YMAX, YOUT, Y0            MAIN 70
C  DOUBLE PRECISION DABS, DBLE, DMAX1                                MAIN 80
      DIMENSION YI(4), ERORI(4), Y0(4,8,4), YDOT(4,4), SAVE(52), RJ(16),MAIN 90
      1                YMAX(4), RW(16), ERO(4), YOUT(4), IPIV(4)      MAIN 100
  
```

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```

COMMON /STAT/ NSTEP,NFE,NJE,NINVS                MAIN 110
DATA LIN/5/,LOUT/6/                             MAIN 120
C*                                                *MAIN 130
C*      SET THE INPUT PARAMETERS                 *MAIN 140
C*                                                *MAIN 150
C*      NP   PROBLEM IDENTIFIER                 *MAIN 160
C*      NC   NUMBER OF TIMES THE SAME PROBLEM IS SOLVED WITH DIFFERENT *MAIN 170
C*      VALUE OF EPS                             *MAIN 180
C*      IP   NUMBER OF TEST POINTS AT WHICH THE SOLUTION IS SOUGHT *MAIN 190
C*      (DOES NOT INCLUDE THE FINAL POINT)      *MAIN 200
C*      N    NUMBER OF DIFFERENTIAL EQUATIONS TO BE SOLVED *MAIN 210
C*      YI   THE INITIAL VALUES OF THE DEPENDENT VARIABLE *MAIN 220
C*      TI   THE INITIAL VALUE OF THE INDEPENDENT VARIABLE *MAIN 230
C*      TF   THE FINAL VALUE OF THE INDEPENDENT VARIABLE *MAIN 240
C*      HI   THE INITIAL STEP-SIZE              *MAIN 250
C*      MF   METHOD FLAG DETERMINING THE MODE OF THE JACOBIAN *MAIN 260
C*      MATRIX EVALUATION                       *MAIN 270
C*      EPS  USERS SPECIFIED ERROR PER STEP    *MAIN 280
C*                                                *MAIN 290
      READ(LIN,10) NP                            MAIN 300
      READ(LIN,10) NC                            MAIN 310
      READ(LIN,10) IP                            MAIN 320
      READ(LIN,10) N                             MAIN 330
10  FORMAT(12)                                   MAIN 340
      READ(LIN,20) (YI(I),I=1,N)                MAIN 350
      READ(LIN,20) TI, TF, HI                  MAIN 360
      DO 1000 K=1,NC                             MAIN 370
      READ(LIN,10) MF                            MAIN 380
      READ(LIN,20) EPS                           MAIN 390
20  FORMAT(3D22.15)                             MAIN 400
      WRITE(LOUT,30) NP, MF, EPS                MAIN 410
30  FORMAT(1H1,10X,13H PROBLEM NO. ,12//        MAIN 420
1    11X,16H METHOD FLAG = ,12//                MAIN 430
2    11X,18H ERROR PER STEP = ,D10.3//         MAIN 440
3    5X,4HTIME,15X,29HS O L U T I O N / E R R O R,26X, MAIN 450
4    7H H-USED,3X,7HNQ-USED,3X,6H STEPS,3X,5H FNS ,3X, MAIN 460
5    5H JACS//)                                MAIN 470
      DO 40 I=1,N                                MAIN 480
C*                                                *MAIN 490
C*      SET THE NORMALIZING VECTOR YMAX         *MAIN 500
C*                                                *MAIN 510
      YMAX(I)=DMAX1(DABS(YI(I)),1.D0)           MAIN 520
C*                                                *MAIN 530
C*      ERORI(I)=0.D0                           MAIN 540
40  Y0(I,1,1)=YI(I)                             MAIN 550
C*                                                *MAIN 560
C*      THE FOLLOWING PARAMETERS ARE USED FOR STATISTICAL PURPOSES *MAIN 570
C*                                                *MAIN 580
C*      NSTEP  NUMBER OF ACCEPTED STEPS         *MAIN 590
C*      NFE    NUMBER OF FUNCTION EVALUATIONS   *MAIN 600
C*      NJE    NUMBER OF JACOBIAN EVALUATIONS   *MAIN 610
C*      NINVS  NUMBER OF CONVERGENCE FACTOR INVERSIONS *MAIN 620
C*                                                *MAIN 630
      NSTEP=0                                    MAIN 640
      NFE=0                                     MAIN 650
      NJE=0                                     MAIN 660
      NINVS=0                                  MAIN 670
C*                                                *MAIN 680
C*      THE FIRST POINT AT WHICH THE SOLUTION IS SOUGHT IS SET BELOW. *MAIN 690
C*      FOR TESTING PURPOSES THE SET OF POINTS TOUT AT WHICH THE *MAIN 700
C*      SOLUTION IS SOUGHT IS DESIGNED SO THAT ITS DENSITY IS LARGE *MAIN 710
C*      AT THE BEGINNING OF THE INTEGRATION INTERVAL AND THEN *MAIN 720
C*      GEOMETRICALLY DECREASES. TOUTP IS THE LAST INTERPOLATION POINT *MAIN 730
C*      HMAX, THE MAXIMUM STEP-SIZE, IS EQUAL TO THE DISTANCE *MAIN 740
C*      BETWEEN THE NEXT AND LAST INTERPOLATION POINTS, MULTIPLIED *MAIN 750
C*      BY 10.                                  *MAIN 760
C*                                                *MAIN 770
      TOUT=TI+(TF-TI)/2.D0**IP                 MAIN 780
      TOUTP=TI                                  MAIN 790
      HMAX=(TOUT-TOUTP)*1.D1                   MAIN 800
C*                                                *MAIN 810
      T =TI                                     MAIN 820
      HNEXT=HI                                  MAIN 830
      HMIN=HI*1.D-2                             MAIN 840
      MAXDER=7                                  MAIN 850
      JSTART=0                                  MAIN 860

```



```

50 CALL STINT (N, T, Y0, YDOT, SAVE, H0, HNEXT, HMIN, HMAX, EPS, MAIN 870
1      YMAX, KFLAG, KNEXT, JSTART, MAXDER, RW, RJ, MF, IPIV) MAIN 880
IF(KFLAG.GE.0) GO TO 200 MAIN 890
WRITE(LOUT,60) KFLAG MAIN 900
60 FORMAT(1H0/20H COMPLETION CODE IS ,I4) MAIN 910
GO TO 800 MAIN 920
200 KFLGP1=KFLAG+1 MAIN 930
C* *MAIN 940
C* CHECK WHETHER THE COMPUTED SOLUTION REACHED BEYOND THE *MAIN 950
C* INTERPOLATION POINT TOUT *MAIN 960
C* *MAIN 970
II=0 MAIN 980
DO 500 I=1,KFLGP1 MAIN 990
II=II+1 MAIN1000
TS=TOUT-T+H0*DBLE(FLOAT(I-1)) MAIN1010
IF(TS.LT.0.D0) GO TO 500 MAIN1020
IF(I-1) 50, 50, 510 MAIN1030
500 CONTINUE MAIN1040
C* *MAIN1050
C* THE SOLUTION REACHED BEYOND TOUT. *MAIN1060
C* PERFORM INTERPOLATION AT TOUT. *MAIN1070
C* *MAIN1080
510 IND=KFLAG+3-II MAIN1090
IF(II.EQ.2) IND=1 MAIN1100
S=(TS-H0)/H0 MAIN1110
DO 600 I=1,N MAIN1120
D=1.D0 MAIN1130
YOUT(I)=Y0(I,1,IND) MAIN1140
DO 600 J=1,JSTART MAIN1150
D=D*(DBLE(FLOAT(J-1))+S)/DBLE(FLOAT(J)) MAIN1160
600 YOUT(I)=YOUT(I)+D*Y0(I,J+1,IND) MAIN1170
CALL ERROR(TOUT,YOUT,N,ERO) MAIN1180
C* *MAIN1190
C* DETERMINE THE MAXIMUM ERROR OF THE INTERPOLATED SOLUTION *MAIN1200
C* COMMITTED SO FAR AND PRINT IT OUT WITH THE SOLUTION AT TOUT *MAIN1210
C* *MAIN1220
DO 610 I=1,N MAIN1230
610 ERORI(I)=DMAX1(ERORI(I),DABS(ERO(I))) MAIN1240
WRITE(LOUT,620) TOUT, (YOUT(I),I=1,N), H0, JSTART, NSTEP, NFE, MAIN1250
1 NJE, (ERO(I),I=1,N) MAIN1260
620 FORMAT( 4H T =,D11.4,4D15.6,D12.3,3X,I4,6X,I4,4X,I4,4X,I4/15X, MAIN1270
1 4D15.6/) MAIN1280
C* *MAIN1290
C* SET THE NEW INTERPOLATION POINT AND CHECK WHETHER AN *MAIN1300
C* INTERPOLATION CAN BE PERFORMED. SET NEW HMAX *MAIN1310
C* *MAIN1320
TOUTP=TOUT MAIN1330
TOUT=TI+(TOUT-TI)*2.D0 MAIN1340
IF(TOUT.GT.TF) GO TO 800 MAIN1350
IF(TOUT.LE.T) GO TO 200 MAIN1360
HMAX=(TOUT-TOUTP)*1.D1 MAIN1370
GO TO 50 MAIN1380
C* *MAIN1390
C* REGARDLESS OF WHETHER OR NOT THE PROBLEM HAS BEEN SOLVED, *MAIN1400
C* PRINT OUT THE FINAL STATISTICS *MAIN1410
C* *MAIN1420
800 WRITE(LOUT,810) MAIN1430
810 FORMAT(1H0/20X,20HINTERPOLATION ERRORS) MAIN1440
WRITE(LOUT,820) (ERORI(I),I=1,N) MAIN1450
820 FORMAT(1H0/10X,23H MAXIMUM ABSOLUTE ERROR//11X,4D22.8) MAIN1460
DO 830 I=1,N MAIN1470
830 ERORI(I)=ERORI(I)/YMAX(I) MAIN1480
WRITE(LOUT,840) (ERORI(I),I=1,N) MAIN1490
840 FORMAT(1H0/10X,23H MAXIMUM RELATIVE ERROR//11X,4D22.8) MAIN1500
WRITE(LOUT,850) NSTEP, NFE, NJE, NINVS MAIN1510
850 FORMAT(1H0/22H PROBLEM COMPLETED IN ,I5,6H STEPS/
1 22X,I5,21H FUNCTION EVALUATIONS/ MAIN1530
2 22X,I5,21H JACOBIAN EVALUATIONS/ MAIN1540
3 22X,I5,18H LU DECOMPOSITIONS) MAIN1550
1000 CONTINUE MAIN1560
STOP MAIN1570
END MAIN1580

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SUBROUTINE STINT (N, T, Y, DY, SAVE, H, HNEXT, HMIN, HMAX, EPS, STNT 10
+ YMAX, KFLAG, KNEXT, JSTART, MAXORD, RW, RJ, MF, STNT 20
+ IPIV) STNT 30
C STNT 40
C*****STNT 50
C* *STNT 60
C* THIS PROGRAM INTEGRATES A SET OF N FIRST ORDER ORDINARY *STNT 70
C* DIFFERENTIAL EQUATIONS. A BLOCK OF THREE OR FOUR SOLUTION POINTS, *STNT 80
C* EACH SEPARATED BY A STEP-SIZE H, IS COMPUTED AT EACH CALL. THE *STNT 90
C* STEP-SIZE MAGNITUDE MAY BE SPECIFIED BY THE USER AT EACH CALL. *STNT 100
C* ALTERNATIVELY, IT MAY BE INCREASED OR DECREASED BY STINT WITHIN *STNT 110
C* THE RANGE DABS(HMIN) TO DABS(HMAX), *STNT 120
C* IN ORDER TO ACHIEVE AS LARGE A STEP AS POSSIBLE, WHILE NOT *STNT 130
C* COMMITTING A SINGLE STEP ERROR WHICH IS LARGER THAN EPS IN THE *STNT 140
C* RMS NORM, WHERE EACH COMPONENT OF THE ERROR VECTOR IS DIVIDED BY *STNT 150
C* THE CORRESPONDING COMPONENT OF YMAX. *STNT 160
C* *STNT 170
C* THE PROGRAM REQUIRES 4 SUBROUTINES NAMED *STNT 180
C* *STNT 190
C* DIFFUN(N, T, Y, DY) *STNT 200
C* PEDERV(N, T, Y, RJ, N0) *STNT 210
C* DEC(N, N0, A, IPIV, IER) *STNT 220
C* SOL(N, N0, A, B, IPIV) *STNT 230
C* *STNT 240
C* DIFFUN EVALUATES THE DERIVATIVES OF THE DEPENDENT VARIABLE Y(I) *STNT 250
C* FOR I=1,...,N AND T AND STORES THE RESULT IN THE ARRAY DY. *STNT 260
C* *STNT 270
C* PEDERV COMPUTES THE PARTIAL DERIVATIVES OF THE DIFFERENTIAL *STNT 280
C* EQUATIONS AT THE VALUES Y(I) FOR I=1,...,N AND T AND STORES THE *STNT 290
C* RESULT IN ARRAY RJ. THUS, RJ(J+(K-1)*N0) IS THE PARTIAL OF DY(J) *STNT 300
C* WITH RESPECT TO Y(K) FOR J,K=1,...,N EVALUATED AT Y(I) FOR I=1,.. *STNT 310
C* ..,N AND T. NOTE--N0 IS THE VALUE OF N ON THE FIRST CALL. IF THE *STNT 320
C* ANALYTIC EXPRESSIONS FOR THE PARTIAL DERIVATIVES ARE NOT *STNT 330
C* AVAILABLE, THEIR APPROXIMATE VALUES CAN BE OBTAINED BY NUMERICAL *STNT 340
C* DIFFERENCING. (SEE PARAMETER MF.) IN THIS CASE SUBROUTINE PEDERV *STNT 350
C* MAY SIMPLY BE *STNT 360
C* *STNT 370
C* SUBROUTINE PEDERV(N, T, Y, RJ, N0) *STNT 380
C* RETURN *STNT 390
C* END *STNT 400
C* *STNT 410
C* DEC PERFORMS AN LU DECOMPOSITION OF A MATRIX A. IF THE *STNT 420
C* DECOMPOSITION IS SUCCESSFUL IER SHOULD BE SET TO 0, OTHERWISE IT *STNT 430
C* SHOULD BE SET TO +1. *STNT 440
C* *STNT 450
C* SOL SOLVES THE LINEAR ALGEBRAIC SYSTEM A*X=B, FOR WHICH THE *STNT 460
C* MATRIX A WAS PROCESSED BY DEC. *STNT 470
C* *STNT 480
C* THIS PROGRAM USES DOUBLE PRECISION FOR ALL FLOATING POINT *STNT 490
C* VARIABLES, EXCEPT FOR THOSE STARTING WITH P, WHICH ARE IN SINGLE *STNT 500
C* PRECISION. *STNT 510
C* *STNT 520
C* TEMPORARY STORAGE SPACE IS PROVIDED (BY THE USER) IN THE *STNT 530
C* ARRAYS IPIV, RJ, RW, AND SAVE. THE ARRAY IPIV HOLDS A VECTOR OF *STNT 540
C* THE SAME NAME. THE ARRAYS RJ AND RW ARE USED TO HOLD MATRICES OF *STNT 550
C* THE SAME NAMES. THE ARRAY SAVE IS PARTITIONED AS FOLLOWS *STNT 560
C* *STNT 570
C* SAVE(J,I) 1.LE.J.LE.8 AND 1.LE.I.LE.N IS USED TO SAVE *STNT 580
C* Y(I,J) IN CASE A STEP (AND HENCE THE WHOLE CYCLE) *STNT 590
C* HAS TO BE REPEATED. *STNT 600
C* SAVE(9,I) 1.LE.I.LE.N IS USED TO STORE THE DERIVATIVE OF THE *STNT 610
C* I-TH DEPENDENT VARIABLE SCALED BY H. *STNT 620
C* SAVE(N9+I,1) IS USED TO STORE THE DERIVATIVES AS THEY ARE *STNT 630
C* COMPUTED BY DIFFUN FOR THE CORRECTOR. IT IS ALSO *STNT 640
C* ACCESSED AS A COMPLETE ARRAY SAVE(N9P1,1). *STNT 650
C* IN ADDITION IT IS USED IN THE ERRCR CONTROL TEST. *STNT 660
C* (N9 = N0 * 9 AND N9P1 = N9 + 1.) *STNT 670
C* SAVE(N10+I,1) IS USED TO HOLD THE CORRECTION TERMS FOR THE ENTIRE *STNT 680
C* CORRECTOR ITERATION IN THE CASE, THE CORRECTOR HAS *STNT 690
C* TO BE REPEATED. (N10 = N0 * 10.) *STNT 700
C* SAVE(N11+I,1) IS USED TO HOLD THE DERIVATIVES EVALUATED AT *STNT 710
C* Y(I)+D AND T, WHERE D IS THE INCREMENT USED IN THE *STNT 720
C* NUMERICAL DIFFERENCING SCHEME INVOKED, IN ORDER TO *STNT 730
C* OBTAIN APPROXIMATE VALUES OF THE PARTIAL *STNT 740
C* DERIVATIVES. (N11 = N0 * 11.) *STNT 750
C* SAVE(N12+I,1) IS USED TO HOLD THE DERIVATIVES EVALUATED AT Y(I) *STNT 760

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C*          AND T IN ORDER TO OBTAIN APPROXIMATE VALUES OF THE *STNT 770
C*          PARTIAL DERIVATIVES. (N12 = N0 * 12.) *STNT 780
C* *STNT 790
C* THE PARAMETERS IN THE CALLING SEQUENCE HAVE THE FOLLOWING MEANING *STNT 800
C* *STNT 810
C* N      THE NUMBER OF FIRST ORDER DIFFERENTIAL EQUATIONS TO BE *STNT 820
C*        INTEGRATED. N MAY BE DECREASED ON SUBSEQUENT CALLS IF *STNT 830
C*        THE NUMBER OF ACTIVE EQUATIONS DECREASES, WITH THE *STNT 840
C*        FIRST ONES BEING THOSE RETAINED. BUT, IT MUST *STNT 850
C*        NOT BE INCREASED WITHOUT CALLING WITH JSTART = 0. *STNT 860
C* T      THE INDEPENDENT VARIABLE. ON ENTRY T IS THE CURRENT *STNT 870
C*        SETTING OF THE INDEPENDENT VARIABLE. ON RETURN TO THE *STNT 880
C*        CALLING PROGRAM, T CORRESPONDS TO THE SETTING OF THE *STNT 890
C*        INDEPENDENT VARIABLE FOR THE MOST FORWARD POINT OBTAINED *STNT 900
C*        THUS FAR. *STNT 910
C* Y      AN N BY 8 BY 4 ARRAY CONTAINING THE DEPENDENT VARIABLES *STNT 920
C*        AND THEIR BACKWARD DIFFERENCES. ON EACH CALL UP TO FOUR *STNT 930
C*        SOLUTION POINTS ARE OBTAINED. THE MOST FORWARD POINT IS *STNT 940
C*        ALWAYS AT Y(I,1,1). THE POINT NEXT TO THE MOST FORWARD *STNT 950
C*        POINT IS RETURNED IN Y(I,1,KFLAG). THE MOST BACKWARD *STNT 960
C*        POINT IN THE NEW BLOCK IS RETURNED IN Y(I,1,2). ONLY THE *STNT 970
C*        INITIAL SOLUTION VALUES, ENTERED IN Y(I,1,1) FOR *STNT 980
C*        I=1,...,N, NEED TO BE SUPPLIED ON THE FIRST *STNT 990
C*        CALL (JSTART = 0). *STNT1000
C*        Y(I,J+1,K) CONTAINS THE J-TH BACKWARD DIFFERENCE OF THE *STNT1010
C*        I-TH DEPENDENT VARIABLE (FOR K=1,...,KFLAG). *STNT1020
C*        IF IT IS DESIRED TO INTERPOLATE TO NON-MESH POINTS, *STNT1030
C*        THESE VALUES CAN BE USED. IF THE CURRENT STEP-SIZE IS *STNT1040
C*        H AND THE VALUE AT T-E (0.LT.DABS(E).LT.DABS(H)) IS *STNT1050
C*        NEEDED, FORM S=E/H AND THEN COMPUTE *STNT1060
C*          NQ *STNT1070
C*          Y(I)(T-E) = SUM Y(I,J+1,K)*B(J) *STNT1080
C*          J=0 *STNT1090
C*        WHERE K, WHICH CORRESPONDS TO T, IS THE FIRST MESH POINT *STNT1100
C*        BEYOND THE POINT T-E, AND B(J+1) = B(J)*(J-S)/(J+1) *STNT1110
C*        WITH B(0) = 1. *STNT1120
C* DY     AN N BY 4 ARRAY. DY(I,K), 1.LE.I.LE.N, 1.LE.K.LE.KFLAG, *STNT1130
C*        CONTAINS THE DERIVATIVE OF THE I-TH DEPENDENT VARIABLE *STNT1140
C*        SCALED BY H. THE DY(I,1) ARRAY NEED NOT BE SUPPLIED AT *STNT1150
C*        THE FIRST CALL (JSTART = 0). *STNT1160
C* SAVE  A BLOCK OF AT LEAST 13*N0 DOUBLE PRECISION FLOATING POINT *STNT1170
C*        LOCATIONS. *STNT1180
C* H     THE STEP-SIZE USED FOR THE JUST COMPLETED BLOCK. *STNT1190
C* HNEXT THE STEP-SIZE FOR THE NEXT BLOCK. ON THE *STNT1200
C*        FIRST CALL (JSTART=0) THE USER MUST SPECIFY AN INITIAL *STNT1210
C*        STEP-SIZE. A GOOD ESTIMATE OF ITS MAGNITUDE IS GIVEN *STNT1220
C*        BY 0.2*(EPS/DABS(E))*0.5, WHERE E IS THE LARGEST *STNT1230
C*        EIGENVALUE OF THE JACOBIAN EVALUATED AT THE INITIAL *STNT1240
C*        VALUES OF T AND Y. THE SIGN OF THE INITIAL HNEXT *STNT1250
C*        SHOULD BE POSITIVE (NEGATIVE) IF THE FINAL TIME IS *STNT1260
C*        GREATER (LESS) THAN THE INITIAL TIME. IF THE INITIAL *STNT1270
C*        STEP-SIZE CHOICE DOES NOT CAUSE AN ERROR GREATER THAN *STNT1280
C*        EPS, IN THE RMS NORM, IT WILL BE ACCEPTED. OTHERWISE, *STNT1290
C*        ITS MAGNITUDE WILL BE DECREASED UNTIL AN ERROR LESS *STNT1300
C*        THAN EPS IS ACHIEVED. THE SUBROUTINE AUTOMATICALLY *STNT1310
C*        ADJUSTS THE STEP-SIZE AFTER THE INITIAL AND SUBSEQUENT *STNT1320
C*        CALLS FOR THE STEP-SIZE OF LARGEST POSSIBLE MAGNITUDE. *STNT1330
C*        THE MAGNITUDE MAY BE ADJUSTED DOWN ON ANY SUBSEQUENT *STNT1340
C*        CALL. NOTE--A MAGNITUDE ADJUSTMENT UP OR A SIGN CHANGE *STNT1350
C*        WILL BE IGNORED. *STNT1360
C* HMIN  DABS(HMIN) IS THE MINIMUM STEP-SIZE MAGNITUDE TO BE *STNT1370
C*        ALLOWED FOR THE NEXT INTEGRATION CYCLE. (ON THE FIRST *STNT1380
C*        CALL (JSTART=0), DABS(HMIN) SHOULD BE CHOSEN SIGNIF- *STNT1390
C*        ICANTLY SMALLER THAN DABS(HNEXT) SO AS TO AVOID START-UP *STNT1400
C*        PROBLEMS IF THE ERROR CRITERION IS NOT MET WITH THE USER *STNT1410
C*        SPECIFIED HNEXT.) NOTE--THE SIGN OF HMIN IS IGNORED. *STNT1420
C*        HMIN MAY BE CHANGED ON SUBSEQUENT CALLS. *STNT1430
C* HMAX  DABS(HMAX) IS THE MAXIMUM STEP-SIZE MAGNITUDE TO BE *STNT1440
C*        ALLOWED FOR THE NEXT INTEGRATION CYCLE. NOTE--IF *STNT1450
C*        DABS(HMAX) IS LESS THAN DABS(HMIN), THEN THE SUBROUTINE *STNT1460
C*        FUNCTIONS AS THOUGH DABS(HMAX) EQUALS DABS(HMIN). HMAX *STNT1470
C*        MAY BE CHANGED ON SUBSEQUENT CALLS. *STNT1480
C* EPS   THE ERROR TEST CONSTANT. THE SINGLE STEP ERROR ESTIMATE *STNT1490
C*        FOR Y, COMPUTED AS A WEIGHTED RMS NORM, MUST NOT EXCEED *STNT1500
C*        EPS. THE WEIGHT FOR THE I-TH ELEMENT IN THE ERROR *STNT1510
C*        ESTIMATE IS 1/YMAX(I). (SFE PARAMETER YMAX.) THE *STNT1520
C*        STEP-SIZE AND/OR ORDER ARE ADJUSTED TO ACHIEVE THIS. *STNT1530

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C          DATA UROUND/4.D-16/, SQRTUR/2.D-8/
C          STNT2310
C          STNT2320
C          STNT2330
C*****STNT2340
C*   THE ARRAY INDEX HOLDS POINTERS AND CONSTANTS FOR THE VARIOUS
C*   ORDER METHODS.  FOR NQ=1,...,7 THE ENTRIES ARE AS FOLLOWS
C*   INDEX(NQ,1)  BASE INDEX FOR B ARRAY (H*DY PREDICTOR).
C*   INDEX(NQ,2)  BASE INDEX FOR C ARRAY (CORRECTOR).
C*****STNT2380
C          STNT2390
C          STNT2400
C          DATA INDEX/ 1, 2, 4, 11, 20, 38, 59,
+          1, 2, 3, 5, 7, 10, 14/
C          STNT2410
C          STNT2420
C          STNT2430
C*****STNT2440
C*   THE COEFFICIENTS APPEARING IN THE NEXT DATA STATEMENTS FOR THE
C*   B ARRAY SHOULD BE DEFINED TO THE MAXIMUM ACCURACY PERMITTED BY
C*   THE MACHINE.  THEY ARE, IN THE ORDER SPECIFIED,...
C*   1
C*   -2, 3
C*   -9/2, -5/4, 11/2
C*   -15/2, -3/4, 13/2, 2
C*   -22/3, -8/3, -17/18, 25/3
C*   -9, -9/4, -7/8, 35/4, 5/4
C*   -125/12, -101/24, -71/36, -37/48, 137/12
C*   -253/24, -1001/240, -707/360, -123/160, 1373/120, 1/10
C*   -57/4, -25/8, -17/10, -169/240, 61/5, -1/20, 31/10
C*   -137/10, -117/20, -46/15, -191/120, -197/300, 147/10
C*   -353/25, -571/100, -1819/600, -79/50, -3919/6000, 1477/100, 7/20
C*   -3529/200, -1889/400, -1609/600, -3527/2400, -931/1500, 3079/200,
C*   -3/20, 17/5
C*   -343/20, -303/40, -253/60, -589/240, -101/75, -23/40, 363/20
C*   -266/15, -221/30, -749/180, -73/30, -803/600, -103/180, 547/30,
C*   1/2
C*   -1316/75, -1151/150, -3689/900, -121/50, -4003/3000, -257/450,
C*   2737/150, -1/5, 1/2
C*****STNT2650
C          STNT2660
C          STNT2670
C          STNT2680
C          DATA B1/
+          1.000, -2.000, 3.000, -4.500, -1.2500, 5.500, -7.500,
+          -.7500, 6.500, 2.000, -7.33333333333333300,
+          -2.666666666666666700, -.944444444444444400,
+          8.33333333333333300, -9.000, -2.2500, -.87500, 8.7500,
+          1.2500, -10.41666666666666700, -4.20833333333333300,
+          -1.97222222222222200, -.770833333333333300,
+          11.41666666666666700, -10.54166666666666700,
+          -4.17083333333333300, -1.96388888888888900, -.7687500,
+          11.44166666666666700, 0.100, -14.2500, -3.12500, -1.700,
+          -.704166666666666700, 12.200, -.0500, 3.100, -13.700,
+          -5.8500, -3.06666666666666700, -1.59166666666666700,
+          -.656666666666666700, 14.700, -14.1200, -5.7100,
+          -3.03166666666666700, -1.5800, -.653166666666666700/
C          STNT2690
C          STNT2700
C          STNT2710
C          STNT2720
C          STNT2730
C          STNT2740
C          STNT2750
C          STNT2760
C          STNT2770
C          STNT2780
C          STNT2790
C          STNT2800
C          STNT2810
C          STNT2820
C          DATA B2/
+          14.7700, 0.3500, -17.64500, -4.722500,
+          -2.681666666666666700, -1.46958333333333300,
+          -.620666666666666700, 15.39500, -.1500, 3.400, -17.1500,
+          -7.57500, -4.21666666666666700, -2.45416666666666700,
+          -1.346666666666666700, -.57500, 18.1500,
+          -17.73333333333333300, -7.36666666666666700,
+          -4.16111111111111100, -2.43333333333333300,
+          -1.33833333333333300, -.57222222222222200,
+          18.23333333333333300, 0.500, -17.54666666666666700,
+          -7.67333333333333300, -4.09888888888888900, -2.4200,
+          -1.33433333333333300, -.57111111111111100,
+          18.24666666666666700, -0.200, 0.500/
C          STNT2830
C          STNT2840
C          STNT2850
C          STNT2860
C          STNT2870
C          STNT2880
C          STNT2890
C          STNT2900
C          STNT2910
C          STNT2920
C          STNT2930
C          STNT2940
C          STNT2950
C*****STNT2960
C*   THE COEFFICIENTS APPEARING IN THE NEXT DATA STATEMENTS FOR THE
C*   C ARRAY SHOULD BE DEFINED TO THE MAXIMUM ACCURACY PERMITTED BY
C*   THE MACHINE.  THEY ARE, IN THE ORDER SPECIFIED,...
C*   -1
C*   -2/3
C*   -6/11, -2/3
C*   -12/25, -16/31
C*   -60/137, -600/1373, -100/193
C*   -20/49, -120/293, -75/184, -1200/2299
C*   -140/363, -60/143, -1050/2437
C*****STNT3070

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C          DATA C/ -1.0D0, -.6666666666666667D0, -.545454545454545D0,      STNT3080
+          -.6666666666666667D0, -.48D0, -.51612903225806452D0,      STNT3090
+          -.43795620437956204D0, -.43699927166788056D0,      STNT3100
+          -.51813471502590674D0, -.40816326530612245D0,      STNT3120
+          -.40955631399317406D0, -.40760869565217391D0,      STNT3130
+          -.52196607220530666D0, -.38567493112947659D0,      STNT3140
+          -.41958041958041958D0, -.43085761181780879D0/      STNT3150
C          STNT3160
C*****STNT3170
C*   THE COEFFICIENTS IN THE PERROR ARRAY ARE USED IN THE ERROR TEST, *STNT3180
C*   THE FIRST TIME IT IS PERFORMED, AS WELL AS IN THE STEP-SIZE/ORDER *STNT3190
C*   SELECTION SEGMENT. PERROR(I+1) = 1/D(I+1), I=1,...,7, WHERE *STNT3200
C*   D(I+1) IS THE DISCRETIZATION ERROR CONSTANT CORRESPONDING TO THE *STNT3210
C*   SECOND PASS OF THE INTEGRATION CYCLE OF ORDER I. PERROR(1) AND *STNT3220
C*   PERROR(9) ARE DEFINED SOLELY FOR PROGRAMMING EASE. THEY ARE NOT *STNT3230
C*   USED. *STNT3240
C*****STNT3250
C          DATA PERROR/ 1.0, 1.0, 1.92857, 2.78161, 3.56735,      STNT3260
+          4.29497, 4.9065, 5.6066, 1.0/      STNT3270
C          STNT3280
C          STNT3290
C*****STNT3300
C*   THE COEFFICIENTS IN THE ARRAY PC ARE USED BOTH IN THE CONVERGENCE *STNT3310
C*   AND ERROR TESTS. THEY ARE THE RECIPROCAL VALUES OF THE *STNT3320
C*   DISCRETIZATION ERROR CONSTANTS FOR EQUATIONS CONSTITUTING THE *STNT3330
C*   METHODS OF ORDER 1 THRU 7. *STNT3340
C*****STNT3350
C          DATA PC/ 2.0, 4.5, 7.3333, 6.0, 10.4167, 9.3, 13.7, 13.8687,      STNT3360
+          9.6904, 17.15, 16.9504, 17.4349, 9.472, 20.7429, 15.921, STNT3370
+          14.7809/      STNT3380
C          STNT3390
C          STNT3400
C*****STNT3410
C*   THE COEFFICIENTS APPEARING IN THE ARRAY PD ARE USED IN THE TESTING *STNT3420
C*   OF THE *OUTDATEDNESS* OF THE ARRAY RW. THE PD(I) ELEMENT CONTAINS *STNT3430
C*   THE AVERAGE VALUE OF THE COEFFICIENTS IN ARRAY C CORRESPONDING TO *STNT3440
C*   ORDER I. *STNT3450
C*****STNT3460
C          DATA PD/ 1.0, 0.6667, 0.6061, 0.4981, 0.4644, 0.4368, 0.412/      STNT3470
C          STNT3480
C          STNT3490
C*****STNT3500
C*****STNT3510
C          STNT3520
C          KFLAG = 0      STNT3530
C          IFAIL = 0      STNT3540
C          IF(JSTART.NE.0) GO TO 80      STNT3550
C*          *STNT3560
C*          INITIALIZATION --- FIRST CALL      *STNT3570
C*          *STNT3580
20 N0 = N      STNT3590
   FN=DBLE(FLOAT(N))      STNT3600
   N9 = N0 * 9      STNT3610
   N9P1 = N9 + 1      STNT3620
   N10 = N9 + N0      STNT3630
   N11 = N10 + N0      STNT3640
   N11P1 = N11 + 1      STNT3650
   N12 = N11 + N0      STNT3660
   N12P1 = N12 + 1      STNT3670
   NSQ = N0*N0      STNT3680
   IWEVAL=+1      STNT3690
   TDL=T      STNT3700
   H=DMAX1(DABS(HMIN), DMIN1(DABS(HNEXT), DABS(HMAX)))      STNT3710
   IF(HNEXT.LT.0.D0) H=-H      STNT3720
30 NQOLD= 0      STNT3730
   ISW1=0      STNT3740
   ISW2=0      STNT3750
   CRATE=1.D0      STNT3760
   CALL DIFFUN (N, T, Y, DY)      STNT3770
   NFE=NFE+1      STNT3780
   DO 40 I=1,N      STNT3790
40   DY(I,1) = H*DY(I,1)      STNT3800
   NQ = 1      STNT3810
   IST = 1      STNT3820
   IDEL = 0      STNT3830
   GO TO 180      STNT3840

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C*
C*          CONTINUE WITH THE STEP-SIZE H
C*
      80 HNEW=DMAX1 (DABS (HMIN) , DMIN1 (DABS (HNEW) , DABS (HNEXT) , DABS (HMAX)))
      IF (H.I.T..0.D0) HNEW=-HNEW
      IF (H.NE.HNEW) GO TO 100
      IST = NQP1
      IDEL = NQP1
      ISW2=0
      GO TO 180

C*
C*          NEW STEP-SIZE --- INTERPOLATE FOR NEW POINTS
C*
      100 RATIO=HNEW/H
      H = RATIO*H
      RC=RC*RATIO*DBLE (PD (NQ) / PDOLD)
      PDOLD = PD (NQ)
      IF (NQ.EQ.1) GO TO 150
      IEQ = NEQ
      D = 0.D0
      DO 140 J=2,NQ
      D = D + RATIO
      IF (D.GT.DBLE (FLOAT (NEQ+NQP1-IEQ))) IEQ=2
      D1 = DBLE (FLOAT (NEQ-IEQ-1)) - D
      DO 140 I=1,N
      D2 = 1.0D0
      D3 = 0.D0
      DO 120 J1=2,NQP1
      D2 = D2*(DBLE (FLOAT (J1)) + D1) / DBLE (FLOAT (J1-1))
      120 D3 = D3 + D2*Y (I, J1, IEQ)
      140 Y (I, J, 1) = D3 + Y (I, 1, IEQ)
      150 IST = NQ
      IDEL = 0
      DO 160 I=1,N
      160 DY (I, 1) = DY (I, 1)*RATIO
      IRET = 3
      GO TO 4000

C*
C*          INITIALIZE SAVE ARRAY
C*
      180 DO 200 I=1,N
      SAVE (9, I) = DY (I, 1)
      DO 200 J=1, IST
      200 SAVE (J, I) = Y (I, J, 1)
      NQST = NQ
      RATIO = 1.0D0
      TOLD = T
      HOLD = H
      220 IF ((NQ.EQ.NQOLD).AND.(FN.EQ.DBLE (FLOAT (N)))) GO TO 300
      IF ((NQ.EQ.NQOLD).AND.(FN.NE.DBLE (FLOAT (N)))) GO TO 280
      IF (MAXORD.I.T.8) GO TO 260
      KFLAG = -7
      GO TO 4185

C*
C*          SET APPROPRIATE PARAMETERS AND CONSTANTS
C*          FOR NEW CYCLE OF ORDER NQ
C*
      260 INDB = INDEX (NQ, 1)
      INDC = INDEX (NQ, 2)
      NEQ = 3+NQ/5
      JSTART = NQ
      NQOLD = NQ
      NQM1=NQ-1
      NQP1 = NQ+1
      Q = DBLE (FLOAT (NQ))
      ENQDWN = 0.5D0/Q
      ENQSAM = 0.5D0 / (Q+1.0D0)
      ENQUP = 0.5D0 / (Q+2.0D0)
      280 FN = DBLE (FLOAT (N))
      EDOWN = FN*(DBLE (PERROR (NQ))*EPS)**2
      EUP = FN*(DBLE (PERROR (NQ+2))*EPS)**2
      ES = FN*(DBLE (PERROR (NQP1))*EPS)**2
      IF (EDOWN.GT.0.D0) GO TO 300
      KFLAG = -5
      GO TO 4185

C*
C*          *STNT3850
C*          *STNT3860
C*          *STNT3870
      STNT3880
      STNT3890
      STNT3900
      STNT3910
      STNT3920
      STNT3930
      STNT3940
      *STNT3950
      *STNT3960
      *STNT3970
      STNT3980
      STNT3990
      STNT4000
      STNT4010
      STNT4020
      STNT4030
      STNT4040
      STNT4050
      STNT4060
      STNT4070
      STNT4080
      STNT4090
      STNT4100
      STNT4110
      STNT4120
      STNT4130
      STNT4140
      STNT4150
      STNT4160
      STNT4170
      STNT4180
      STNT4190
      STNT4200
      STNT4210
      *STNT4220
      *STNT4230
      *STNT4240
      STNT4250
      STNT4260
      STNT4270
      STNT4280
      STNT4290
      STNT4300
      STNT4310
      STNT4320
      STNT4330
      STNT4340
      STNT4350
      STNT4360
      STNT4370
      *STNT4380
      *STNT4390
      *STNT4400
      *STNT4410
      STNT4420
      STNT4430
      STNT4440
      STNT4450
      STNT4460
      STNT4470
      STNT4480
      STNT4490
      STNT4500
      STNT4510
      STNT4520
      STNT4530
      STNT4540
      STNT4550
      STNT4560
      STNT4570
      STNT4580
      STNT4590
      *STNT4600

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C*          CHECK FOR REEVALUATION OF JACOBIAN
C*
300 IF(IWEVAL.GT.0) GO TO 320
IF(DABS(RC-1.D0).GE.4.D-1) IWEVAL = +2
IF((TDL-TOLD)*H.GT.0.D0) GO TO 320
IF(DABS(RC-1.D0).GE.8.D-1) IWEVAL = +1
320 INDBB = INDB
INDCC = INDC
DO 1000 IEQ=1,NEQ
  IST = MOD (IEQ, NEQ) + 1
  T = T + H
  BND=FN*(DBLE(PC(INDCC))*ENQUP*EPS)**2
  E=ES
  IF(IEQ.GT.2) E=FN*(DBLE(PC(INDCC))*EPS)**2
C*
C*          PREDICT Y AND DY FOR THE NEXT MESH POINT
C*
DO 460 I=1,N
  D = DY(I,IEQ)
  D1 = Y(I,1,IEQ) + Q*D
  D2 = B(INDBB+NQM1)*D
  IF (NQ-2) 440, 400, 340
340 IF (IEQ-3) 400, 380, 360
360 D2 = D2 + B(INDBB+NQP1)*DY(I,3)
380 D2 = D2 + B(INDBB+NQ)*DY(I,2)
400 DO 420 J=2,NQ
  D = Y(I,J,IEQ)
  D3 = DBLE(FLOAT(J-1))
  D1 = D1 + D*(D3-Q)/D3
420 D2 = D2 + D*B(INDBB+J-2)
440 Y(I,1,IST) = D1
460 DY(I,IST) = D2
IF(NQ.LE.2) GO TO 470
IF(IEQ.GT.1) INDBB = INDBB + NQ + IEQ - 2
C*
C*          ITERATE THE CORRECTOR UP TO THREE TIMES. ACCUMULATE THE
C*          CORRECTION TERMS IN SAVE(N10+I,1) FOR REDOING THE ENTIRE
C*          CORRECTOR IF CONVERGENCE IS NOT ACHIEVED
C*
470 D1 = C(INDCC)
500 DO 510 I=1,N
510 SAVE(N10+I,1) = 0.D0
DO 780 J=1,3
  CALL DIFFUN (N, T, Y(1,1,IST), SAVE(N9P1,1))
  NFE=NFE+1
  IF(IWEVAL-1) 680, 520, 620
520 IND = 1
  IF(IEQ.EQ.2) IND = 1
  GO TO (610, 530), MF
530 CALL DIFFUN (N, T-H*DBLE(FLOAT(1+IEQ-IND)), Y(1,1,IND),
+   SAVE(N12P1,1))
  NFE=NFE+1
  D = 0.D0
  DO 540 I=1,N
540 D = D + SAVE(N12+I,1)**2
  D = DABS(H)*1.D3*UROUND*DSQRT(D)
  NJE = NJE + 1
  J0 = 0
  DO 560 J1=1,N
  DI = SQRTUR*YMAX(J1)
  DI = DMAX1(DI,D)
  YJ1 = Y(J1,1,IND)
  Y(J1,1,IND) = Y(J1,1,IND) + DI
  CALL DIFFUN (N, T-H*DBLE(FLOAT(1+IEQ-IND)), Y(1,1,IND),
+   SAVE(N11P1,1))
  NFE = NFE + 1
  DO 550 I=1,N
550 RJ(I+J0) = (SAVE(N11+I,1) - SAVE(N12+I,1))/DI
  J0 = J0 + N0
  Y(J1,1,IND) = YJ1
560 CONTINUE
610 GO TO 620
CALL PEDERN(N, T-H*DBLE(FLOAT(1+IEQ-IND)), Y(1,1,IND), RJ, N0)
NJE=NJE+1
620 D = D1*H
DO 640 I=1,NSQ

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*STNT4610
*STNT4620
STNT4630
STNT4640
STNT4650
STNT4660
STNT4670
STNT4680
STNT4690
STNT4700
STNT4710
STNT4720
STNT4730
STNT4740
*STNT4750
*STNT4760
*STNT4770
STNT4780
STNT4790
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STNT4900
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STNT4930
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*STNT4950
*STNT4960
*STNT4970
*STNT4980
*STNT4990
STNT5000
STNT5010
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STNT5290
STNT5300
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STNT5340
STNT5350
STNT5360

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640      RW(I) = RJ(I)*D          STNT5370
      J1 = 1                    STNT5380
      DO 660 I=1,N              STNT5390
        RW(J1) = 1.D0 + RW(J1) STNT5400
660      J1 = J1 + N0 + 1       STNT5410
      CALL DEC (N, N0, RW, IPIV, IER) STNT5420
      NINVS=NINVS+1            STNT5430
      IWEVAL = -IEQ            STNT5440
      RC=1.D0                  STNT5450
      PDOLD=PD(NQ)             STNT5460
      IF (IER.NE.0) GO TO 800    STNT5470
680      DO 700 I=1,N          STNT5480
700      SAVE(N9+I,1) = DY(I,IST) - H*SAVE(N9+I,1) STNT5490
      CALL SOL (N, N0, RW, SAVE(N9P1,1), IPIV) STNT5500
      D2 = 0.D0                STNT5510
      J3 = N9                  STNT5520
      DO 760 I=1,N             STNT5530
        J3 = J3 + 1            STNT5540
        J2 = J3 + N0           STNT5550
        SAVE(J2,1) = SAVE(J2,1) + SAVE(J3,1) STNT5560
        Y(I,1,IST) = Y(I,1,IST) + D1*SAVE(J3,1) STNT5570
        DY(I,1,IST) = DY(I,1,IST) - SAVE(J3,1) STNT5580
760      D2 = D2 + (SAVE(J3,1)/YMAX(I))*2 STNT5590
      IF (J.NE.1) CRATE = DMAX1(CRATE*9.D-1,D2/D3) STNT5600
      IF (D2*DMIN1(1.D0,2.D0*CRATE).GT.BND) GO TO 770 STNT5610
      GO TO 940                STNT5620
770      D3 = D2                STNT5630
780      CONTINUE              STNT5640
C*                               *STNT5650
C* CORRECTOR FAILED TO CONVERGE IN THREE ITERATIONS. *STNT5660
C* IF JACOBIAN WAS REEVALUATED DURING THIS CYCLE, STEP-SIZE IS *STNT5670
C* REDUCED TO 3/10 OF H. OTHERWISE, JACOBIAN IS REEVALUATED *STNT5680
C*                               *STNT5690
      TDL = TOLD                STNT5700
      IF (IWEVAL.EQ.0) GO TO 900 STNT5710
800      IF (DABS(H).LE.(1.00001D0*DABS(HMIN))) IF (NQ-2) 4180, 1080, 4140 STNT5720
      RATIO = RATIO*0.3D0       STNT5730
      IRET = 1                  STNT5740
      ISW1=0                    STNT5750
      ISW2=1                    STNT5760
      IWEVAL = +2               STNT5770
      GO TO 3000                STNT5780
900      DO 920 I=1,N           STNT5790
        D = SAVE(N10+I,1)       STNT5800
        Y(I,1,IST) = Y(I,1,IST) - D1*D STNT5810
920      DY(I,IST) = DY(I,IST) + D STNT5820
      IWEVAL = +1              STNT5830
      GO TO 500                 STNT5840
C*                               *STNT5850
C* CORRECTOR CONVERGED. THE BACKWARD DIFFERENCES OF ORDER *STNT5860
C* ONE THROUGH NQ AT THE NEW MESH POINT ARE COMPUTED *STNT5870
C*                               *STNT5880
940      DO 980 I=1,N           STNT5890
        DO 960 J=1,NQ           STNT5900
960      Y(I,J+1,IST) = Y(I,J,IST) - Y(I,J,IEQ) STNT5910
      IF (IEQ.EQ.1) GO TO 980    STNT5920
C*                               *STNT5930
C* THE (NQ+1)-ST BACKWARD DIFFERENCE FOR ALL BUT THE FIRST *STNT5940
C* MESH POINT IS ESTABLISHED *STNT5950
C*                               *STNT5960
      SAVE(N9+I,1) = Y(I,NQ1,IST) - Y(I,NQ1,IEQ) STNT5970
980      CONTINUE              STNT5980
      IF (IEQ.EQ.NEQ) GO TO 1015 STNT5990
      IF (NQ.LE.2) GO TO 1010    STNT6000
      IF ((IEQ.GT.1).OR.(NQ.EQ.6)) INDCC = INDCC + 1 STNT6010
1010     IF (IEQ.EQ.1) GO TO 1000 STNT6020
C*                               *STNT6030
C* ERROR TEST FOR ALL BUT THE FIRST MESH POINT IS PERFORMED *STNT6040
C*                               *STNT6050
1015     D = 0.D0              STNT6060
        DO 1020 I=1,N           STNT6070
1020     D = D + (SAVE(N9+I,1)/YMAX(I))*2 STNT6080
        IF (D.GT.E) GO TO 1030  STNT6090
1000     CONTINUE              STNT6100
      GO TO 1160                STNT6110
C*                               *STNT6120

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C*          THE ERROR CRITERION WAS NOT MET
C*
1030 IFAIL=IFAIL+1
      IF(IFAIL.GT.2) GO TO 1080
      IF(DABS(H).LE.(DABS(HMIN)*1.00001D0)) IF(NQ-2) 1140, 4120, 4140
      IWEVAL = +2
      IF(IFAIL.EQ.1) GO TO 1200
      TDL=T
      RATIO = RATIO*5.D-1
      IRET = 1
      ISW1=0
      ISW2=1
      GO TO 3000

C*
C*          START OVER WITH ORDER 1 METHOD
C*
1080 IFAIL = 0
      DO 1090 I=1,N
1090  Y(1,1,1) = SAVE(1,I)
      T = TOLD
      H = H*DMAX1(1.D-1,DABS(HMIN/H))
      IWEVAL = +2
      GO TO 3000

C*
C*          A NEW BLOCK OF SOLUTION POINTS HAS BEEN ACCEPTED
C*
1140 IWEVAL = 0
      NSTEP = NSTEP + IEQ
      KFLAG = -IEQ
      RETURN
1160 IWEVAL = 0
      E = ES
      KFLAG = NEQ
      NSTEP = NSTEP + KFLAG
      HNEW=H

C*
C*          CHECK FOR CONTINUATION WITH THE SAME H AND NQ
C*
      IF(ISW2.EQ.1) GO TO 1420
      IF(NQ.GT.3) ISW1=1-ISW1
      IF(ISW1.EQ.1) GO TO 1420

C*
C*          NEW STEP-SIZE AND/OR ORDER SELECTION
C*
1200 RRSAME=1.2D0*(D/E)**ENQSAM
      IF(IFAIL.NE.0) GO TO 1380
      RMAX = 1.D-4
      DF=DBLE(FLOAT(NEQ+NQM1))
      IF(NQ.NE.1) RMAX=(Q-1.D0)/DF
      RRSAME=DMAX1(RRSAME,RMAX)
      RRUP = 1.D20
      RRDOWN = 1.D20
      IF(NQ.GE.MAXORD) GO TO 1240
      D = 0.0D0
      DO 1220 I=1,N
        D1 = Y(I,NQP1,NEQ) - Y(I,NQP1,NEQ-1)
1220  D = D + ((SAVE(N9+I,1) - D1)/YMAX(I))**2
      RRUP = 1.2D0*(D/EUP)**ENQUP
      RMAX=Q/DF
      RRUP=DMAX1(RRUP,RMAX)
1240 IF(NQ.EQ.1) GO TO 1280
      D = 0.0D0
      DO 1260 I=1,N
1260  D = D + (Y(I,NQP1,1)/YMAX(1))**2
      RRDOWN = 1.2D0*(D/EDOWN)**ENQDWN
      RMAX = 1.D-4
      IF(NQ.NE.2) RMAX=(Q-2.D0)/DF
      RRDOWN=DMAX1(RRDOWN,RMAX)
1280 IF (RRSAME.GT.RRUP) IF (RRUP-RRDOWN) 1340, 1300, 1300
      IF (RRSAME.LE.RRDOWN) GO TO 1320
1300 NEWQ = NQM1
      D=1.D0/RRDOWN
      GO TO 1360
1320 NEWQ = NQ
      D=1.D0/RRSAME
      GO TO 1360

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*STNT6130
*STNT6140
STNT6150
STNT6160
STNT6170
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STNT6190
STNT6200
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*STNT6270
*STNT6280
STNT6290
STNT6300
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STNT6330
STNT6340
STNT6350
*STNT6360
*STNT6370
*STNT6380
STNT6390
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STNT6860
STNT6870
STNT6880

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1340 NEWQ = NQF1
      D=1.D0/RRUP
1360 IF (D-1.D0) 1420, 1400, 1400
1380 RATIO=RATIO/RRSAME
      IRET = 1
      ISW1=0
      ISW2=1
      GO TO 3000
1400 HNEW = H*D
      NQ = NEWQ
1420 DO 1460 I=1,N
      D=YMAX(I)
      DO 1440 J=1,NEQ
1440   D=DMAX1(D,DABS(Y(I,1,J)))
1460   YMAX(I)=D
      HNEXT = HNEW
      KNEXT = 3+NQ/5
      RETURN
C*
C*   THIS SECTION IS USED WHEN STEP-SIZE OR ORDER IS CHANGED
C*   DURING THE CYCLE.  STARTING VALUES ARE RETRIEVED FROM THE
C*   SAVE ARRAY.
C*
3000 RATIO = DMAX1 (DABS(HMIN/HOLD), DMIN1(RATIO, 1.0D0))
      T = TOLD
      IF (RATIO.LT.1.0D0) IF (IDEL) 3030, 3030, 3080
      DO 3020 I=1,N
      DY(I,1) = SAVE(9,I)
      DO 3020 J=1,NQ
3020   Y(I,J,1) = SAVE(J,I)
      GO TO (320,260), IRET
C*
C*   THE (NQST+1)-ST ORDER BACKWARD DIFFERENCE IS ESTABLISHED
C*
3030 IDEL = NQST+1
      IF ((NQST.LT.2) .OR. (NQ.LT.2)) GO TO 3080
      D = DBLE(FLOAT(NQST))
      DO 3060 I=1,N
      D1 = SAVE(9,I)
      DO 3040 J=2,NQST
3040   D1 = D1 - SAVE(J,I)/DBLE(FLOAT(J-1))
3060   SAVE(IDEL,I) = D*D1
C*
C*   INTERPOLATE FOR NEW POINTS
C*
3080 DO 3140 I=1,N
      DY(I,1) = RATIO*SAVE(9,I)
      IF (NQ.LT.2) GO TO 3140
      D1 = 0.0D0
      DO 3120 J=2,NQ
      D1 = D1 + RATIO
      D2 = 1.0D0
      D = 0.0D0
      DO 3100 J1=2, IDEL
      D2 = D2*(DBLE(FLOAT(J1-2)) - D1)/DBLE(FLOAT(J1-1))
3100   D = D + D2*SAVE(J1,I)
3120   Y(I,J,1) = D + SAVE(1,I)
3140   Y(I,1,1) = SAVE(1,I)
      H = HOLD*RATIO
C*
C*   FORM THE BACKWARD DIFFERENCES
C*
4000 IF (NQ.LT.2) GO TO 4040
      NQM1 = NQ-1
      DO 4020 I=1,N
      DO 4020 J=1,NQM1
      J0=J + 1
      DO 4020 J1=J0,NQ
      J2 = NQ-J1+J+1
4020   Y(I,J2,1) = Y(I,J2-1,1) - Y(I,J2,1)
4040 GO TO (320,260,180), IRET
C*
C*   CONVERGENCE OR ERROR PROBLEMS
C*
4120 NQ = 1
      GO TO 4160

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STNT6890
STNT6900
STNT6910
STNT6920
STNT6930
STNT6940
STNT6950
STNT6960
STNT6970
STNT6980
STNT6990
STNT7000
STNT7010
STNT7020
STNT7030
STNT7040
STNT7050
STNT7060
*STNT7070
*STNT7080
*STNT7090
*STNT7100
*STNT7110
STNT7120
STNT7130
STNT7140
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STNT7550
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STNT7570
STNT7580
STNT7590
*STNT7600
*STNT7610
*STNT7620
STNT7630
STNT7640

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4140 NQ = 2                                STNT7650
4160 IFAIL = 0                             STNT7660
      IRET = 2                             STNT7670
      IWEVAL = +2                          STNT7680
      GO TO 30000                           STNT7690
4180 KFLAG = -6                            STNT7700
4185 J1 = NQST + 1                         STNT7710
      DO 4190 I=1,N                        STNT7720
          DY(I,1) = SAVE(9,I)              STNT7730
          DO 4190 J=1,J1                    STNT7740
4190     Y(I,J,1) = SAVE(J,I)             STNT7750
      H = HOLD                             STNT7760
      T = TOLD                             STNT7770
      JSTART = NQST                        STNT7780
      RETURN                               STNT7790
C                                          STNT7800
C*****STNT7810
C*                                          *STNT7820
C*          --- END OF SUBROUTINE STINT --- *STNT7830
C*                                          *STNT7840
C*****STNT7850
C                                          STNT7860
      END                                  STNT7870

      SUBROUTINE DEC (N, NDIM, A, IP, IER)   DEC0 10
C                                          DEC0 20
      INTEGER N,NDIM,IP,IER,NM1,K,KP1,M,I,J DEC0 30
      DOUBLE PRECISION A, T                DEC0 40
C                                          DEC0 50
      DOUBLE PRECISION DABS
      DIMENSION A(NDIM,N), IP(N)          DEC0 60
-----DEC0 70
C MATRIX TRIANGULARIZATION BY GAUSSIAN ELIMINATION. DEC0 80
C INPUT.. DEC0 90
C   N = ORDER OF MATRIX. DEC0 100
C   NDIM = DECLARED DIMENSION OF ARRAY A . DEC0 110
C   A = MATRIX TO BE TRIANGULARIZED. DEC0 120
C OUTPUT.. DEC0 130
C   A(I,J), I.LE.J = UPPER TRIANGULAR FACTOR, U . DEC0 140
C   A(I,J), I.GT.J = MULTIPLIERS = LOWER TRIANGULAR FACTOR, I - L. DEC0 150
C   IP(K), K.LT.N = INDEX OF K-TH PIVOT ROW. DEC0 160
C   IP(N) = (-1)**(NUMBER OF INTERCHANGES) OR 0 . DEC0 170
C   IER = 0 IF MATRIX A IS NONSINGULAR, OR K IF FOUND TO BE DEC0 180
C   SINGULAR AT STAGE K. DEC0 190
C USE SOL TO OBTAIN SOLUTION OF LINEAR SYSTEM. DEC0 200
C DETERM(A) = IP(N)*A(1,1)*A(2,2)*...*A(N,N). DEC0 210
C IF IP(N)=0, A IS SINGULAR, SOL WILL DIVIDE BY ZERO. DEC0 220
C DEC0 230
C REFERENCE.. DEC0 240
C   C. B. MOLER, ALGORITHM 423, LINEAR EQUATION SOLVER, DEC0 250
C   C.A.C.M. 15 (1972), P. 274. DEC0 260
-----DEC0 270
      IER = 0                               DEC0 280
      IP(N) = 1                             DEC0 290
      IF (N .EQ. 1) GO TO 70                 DEC0 300
      NM1 = N - 1                           DEC0 310
      DO 60 K = 1,NM1                        DEC0 320
          KP1 = K + 1                        DEC0 330
          M = K                              DEC0 340
          DO 10 I = KP1,N                    DEC0 350
10             IF (DABS(A(I,K)) .GT. DABS(A(M,K))) M = I DEC0 360
              IP(K) = M                     DEC0 370
              T = A(M,K)                    DEC0 380
              IF (M .EQ. K) GO TO 20         DEC0 390
              IP(N) = -IP(N)                DEC0 400
              A(M,K) = A(K,K)               DEC0 410
              A(K,K) = T                     DEC0 420
20             IF (T .EQ. 0.D0) GO TO 80     DEC0 430
              T = 1.D0/T                    DEC0 440
              DO 30 I = KP1,N                DEC0 450
30                 A(I,K) = -A(I,K)*T      DEC0 460
              DO 50 J = KP1,N                DEC0 470
                  T = A(M,J)                DEC0 480
                  A(M,J) = A(K,J)           DEC0 490
                  A(K,J) = T                DEC0 500
                  IF (T .EQ. 0.D0) GO TO 50 DEC0 510

```

```

      DO 40 I = KP1,N                                DEC0 520
40      A(I,J) = A(I,J) + A(I,K)*T                   DEC0 530
50      CONTINUE                                     DEC0 540
60      CONTINUE                                     DEC0 550
70      K = N                                         DEC0 560
      IF (A(N,N) .EQ. 0.D0) GO TO 80                 DEC0 570
      RETURN                                          DEC0 580
80      IER = K                                       DEC0 590
      IP(N) = 0                                       DEC0 600
      RETURN                                          DEC0 610
C----- END OF SUBROUTINE DEC -----DEC0 620
      END                                            DEC0 630

      SUBROUTINE SOL (N, NDIM, A, B, IP)              SOL0 10
C                                                    SOL0 20
      INTEGER N,NDIM,IP,NM1,K,KP1,M,I,KB,KM1        SOL0 30
      DOUBLE PRECISION A,B,T                         SOL0 40
      DIMENSION A(NDIM,N), B(N), IP(N)             SOL0 50
C----- SOL0 60
C SOLUTION OF LINEAR SYSTEM, A*X = B .             SOL0 70
C INPUT..                                           SOL0 80
C N = ORDER OF MATRIX.                             SOL0 90
C NDIM = DECLARED DIMENSION OF ARRAY A .           SOL0 100
C A = TRIANGULARIZED MATRIX OBTAINED FROM DEC.     SOL0 110
C B = RIGHT HAND SIDE VECTOR.                     SOL0 120
C IP = PIVOT VECTOR OBTAINED FROM DEC.             SOL0 130
C DO NOT USE IF DEC HAS SET IER .NE. 0.           SOL0 140
C OUTPUT..                                         SOL0 150
C B = SOLUTION VECTOR, X .                         SOL0 160
C----- SOL0 170
      IF (N .EQ. 1) GO TO 50                          SOL0 180
      NM1 = N - 1                                    SOL0 190
      DO 20 K = 1,NM1                                 SOL0 200
          KP1 = K + 1                                 SOL0 210
          M = IP(K)                                   SOL0 220
          T = B(M)                                    SOL0 230
          B(M) = B(K)                                 SOL0 240
          B(K) = T                                    SOL0 250
          DO 10 I = KP1,N                             SOL0 260
10             B(I) = B(I) + A(I,K)*T                 SOL0 270
20             CONTINUE                               SOL0 280
          DO 40 KB = 1,NM1                             SOL0 290
              KM1 = N - KB                           SOL0 300
              K = KM1 + 1                             SOL0 310
              B(K) = B(K)/A(K,K)                       SOL0 320
              T = -B(K)                                 SOL0 330
              DO 30 I = 1,KM1                           SOL0 340
30                 B(I) = B(I) + A(I,K)*T             SOL0 350
40                 CONTINUE                           SOL0 360
50                 B(1) = B(1)/A(1,1)                 SOL0 370
              RETURN                                    SOL0 380
C----- END OF SUBROUTINE SOL -----SOL0 390
      END                                            SOL0 400

      SUBROUTINE DIFFUN(N, T, Y, YDOT)              TEST 10
      DOUBLE PRECISION T, T1, T2, W1, W2, W3, W4, S, Y, YDOT TEST 20
      DIMENSION Y(N), YDOT(N)                       TEST 30
      T1=-Y(1)+Y(2)+Y(3)+Y(4)                       TEST 40
      T2=Y(1)-Y(2)+Y(3)+Y(4)                       TEST 50
      W1=5.D-1*(T1*T1-T2*T2)                         TEST 60
      W2=T1*T2                                       TEST 70
      W3=(Y(1)+Y(2)-Y(3)+Y(4))**2                   TEST 80
      W4=(Y(1)+Y(2)+Y(3)-Y(4))**2                   TEST 90
      S=0.D0                                          TEST 100
      DO 10 I=1,3                                    TEST 110
10 S=S-1.D-3*Y(I)                                  TEST 120
      S=S+1.D-3*Y(4)                                  TEST 130
      YDOT(1)=-2.5D-1*(9.8D2*Y(1)+1.02D3*Y(2)-9.8D2*Y(3)+1.02D3*Y(4)-S) TEST 140
      + 1.25D-1*(-W1+W2+W3+W4)                       TEST 150
      YDOT(2)=-2.5D-1*(1.02D3*Y(1)+9.8D2*Y(2)-1.02D3*Y(3)+9.8D2*Y(4)-S) TEST 160
      + 1.25D-1*( W1-W2+W3+W4)                       TEST 170
      YDOT(3)=-2.5D-1*(-1.02D3*Y(1)-9.8D2*Y(2)+9.8D2*Y(3)-1.02D3*Y(4)-S) TEST 180
      + 1.25D-1*( W1+W2-W3+W4)                       TEST 190
      YDOT(4)=-2.5D-1*(9.8D2*Y(1)+1.02D3*Y(2)-1.02D3*Y(3)+9.8D2*Y(4)+S) TEST 200
      + 1.25D-1*( W1+W2+W3-W4)                       TEST 210

```

RETURN	TEST 220
END	TEST 230
SUBROUTINE PEDERV(N, T, Y, P0, N0)	TEST 240
DOUBLE PRECISION P0, Y, T	TEST 250
DIMENSION P0(16), Y(N)	TEST 260
P0(1) = (-9.800001D2+Y(1)+Y(3)+Y(4)+3.D0*Y(2))*2.5D-1	TEST 270
P0(5) = (-1.0200001D3+Y(2)-Y(3)-Y(4)+3.D0*Y(1))*2.5D-1	TEST 280
P0(9) = (9.799999D2+Y(1)-Y(2)-Y(4)+3.D0*Y(3))*2.5D-1	TEST 290
P0(13) = (-1.019999D3+Y(1)-Y(2)-Y(3)+3.D0*Y(4))*2.5D-1	TEST 300
P0(2) = P0(5)	TEST 310
P0(6) = P0(1)	TEST 320
P0(10) = -P0(13)	TEST 330
P0(14) = -P0(9)	TEST 340
P0(3) = P0(10)	TEST 350
P0(7) = P0(9)	TEST 360
P0(11) = P0(1)	TEST 370
P0(15) = -P0(5)	TEST 380
P0(4) = P0(14)	TEST 390
P0(8) = P0(13)	TEST 400
P0(12) = P0(15)	TEST 410
P0(16) = P0(1)	TEST 420
RETURN	TEST 430
END	TEST 440
SUBROUTINE ERROR(T, Y, N, ER0)	TEST 450
DOUBLE PRECISION T, T1, T2, T3, W1, W2, W3, W4, X, Y, ER0	TEST 460
C DOUBLE PRECISION DEXP, DSIN, DCOS, DABS	TEST 470
DIMENSION ER0(N), Y(N)	TEST 480
X = 1.D1*T	TEST 490
IF(X.GT.1.65D2) GO TO 20	TEST 500
T1 = 1.D0+DEXP(-X)*(9.D0*DCOS(X)+1.D1*DSIN(X))	TEST 510
T2 = DEXP(-X)*(1.D1*DCOS(X)-9.D0*DSIN(X))	TEST 520
IF(DABS(T2).LE.1.D-15) GO TO 10	TEST 530
T3 = 1.D1/(T1*T1+T2*T2)	TEST 540
GO TO 11	TEST 550
10 T3 = 1.D1/T1**2	TEST 560
11 W1 = -(T1+T2)*T3	TEST 570
W2 = (T1-T2)*T3	TEST 580
GO TO 21	TEST 590
20 W1 = -1.D1	TEST 600
W2 = 1.D1	TEST 610
21 X = 1.D3*T	TEST 620
IF(X.GT.1.65D2) GO TO 22	TEST 630
W3 = 5.D2/(1.D0-1.001D3*DEXP(X))	TEST 640
GO TO 23	TEST 650
22 W3 = 0.D0	TEST 660
23 W4 = 5.D-4/(1.D0-1.001D0*DEXP(1.D-3*T))	TEST 670
ER0(1) = -W1+W2+W3+W4	TEST 680
ER0(2) = W1-W2+W3+W4	TEST 690
ER0(3) = W1+W2-W3+W4	TEST 700
ER0(4) = W1+W2+W3-W4	TEST 710
DO 30 I = 1, N	TEST 720
30 ER0(I) = ER0(I) - Y(I)	TEST 730
RETURN	TEST 740
END	TEST 750
20 .	TEST 760
6	TEST 770
15	TEST 780
04	TEST 790
0.0000000000000000D+00-2.0000000000000000D+00-1.0000000000000000D+00	TEST 800
-1.0000000000000000D+00	TEST 810
0.0000000000000000D+00 1.0000000000000000D+02 5.0000000000000000D-09	TEST 820
1	TEST 830
1.0000000000000000D-03	TEST 840
1	TEST 850
1.0000000000000000D-05	TEST 860
1	TEST 870
1.0000000000000000D-07	TEST 880
2	TEST 890
1.0000000000000000D-03	TEST 900
2	TEST 910
1.0000000000000000D-05	TEST 920
2	TEST 930
1.0000000000000000D-07	TEST 940

ALGORITHM 535

The QZ Algorithm To Solve the Generalized Eigenvalue Problem for Complex Matrices [F2]

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Key Words and Phrases: eigenvalues, generalized eigenvalue problem
CR Categories: 5.14
Language: Fortran

DESCRIPTION

Three Fortran subroutines are provided that implement a complex form of the QZ algorithm for finding λ and z such that $Az = \lambda Bz$, where A and B are complex N by N matrices. The subroutines are complex analogs of those for the corresponding real problem included in the EISPACK eigensystem package [1]; the original QZ algorithm is described in [3] and [4]. The complex QZ algorithm shares the fundamental property of the real QZ algorithm in being unaffected by singularity or near singularity of B .

Subroutine CQZHES implements the first step of the algorithm wherein A and B are simultaneously reduced by unitary transformations to upper Hessenberg and upper triangular form, respectively. Subroutine CQZVAL implements an iterative process that reduces A to upper triangular form while maintaining the triangular form of B . At the completion of this step, the eigenvalues are derivable from the corresponding diagonal elements of the reduced A and B . Finally, if eigenvectors are desired, subroutine CQZVEC applies the accumulated transformations from the two earlier steps onto the eigenvectors of the triangular problem.

No facility is provided for obtaining just a few eigenvectors; a subroutine employing inverse iteration might be included for this purpose, proceeding from the Hessenberg A , the triangular B , and an approximate eigenvalue of the problem. Inverse iteration could be expected to save time, but it might be somewhat less accurate and reliable, especially in the presence of close eigenvalues.

Also, no facility is included for balancing A and B . Optimally scaled matrices could result in further improvement of accuracy obtainable with unitary transformations, but no automatic scaling technique is known at this time.

The subroutines can be invoked with consecutive statements:

```
CALL CQZHES (NM, N, AR, AI, BR, BI, MATZ, ZR, ZI)
CALL CQZVAL (NM, N, AR, AI, BR, BI, EPS1, ALFR, ALFI, BETA, MATZ, ZR, ZI,
            IERR)
```

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THE FULL MATRIX A OF ORDER 5 IS (PRINTED BY ROWS)									
-238.	-344.I	86.	178.I	164.	240.I	-166.	-308.I	56.	158.I
76.	152.I	-96.	-128.I	40.	-32.I	60.	184.I	-60.	-136.I
118.	284.I	55.	-182.I	-13.	460.I	34.	-192.I	-176.	-214.I
-314.	-160.I	132.	78.I	114.	296.I	-90.	-164.I	-424.	-374.I
-54.	-24.I	-205.	-400.I	109.	148.I	158.	312.I	-38.	-96.I
THE FULL MATRIX B OF ORDER 5 IS (PRINTED BY ROWS)									
388.	94.I	-386.	-122.I	-250.	-14.I	556.	130.I	-396.	-62.I
-304.	-76.I	384.	64.I	-160.	16.I	-240.	-92.I	240.	68.I
-658.	-136.I	-73.	100.I	-109.	-250.I	-118.	100.I	406.	96.I
-640.	-10.I	204.	-42.I	-692.	-90.I	288.	66.I	-192.	154.I
-162.	-72.I	631.	158.I	131.	52.I	-758.	-184.I	278.	76.I
ALFR(I)			ALFI(I)			BETA(I)			
1	3.299829E 02		4.061311E 02		4.315166E 02				
2	-1.368646E 02		-1.824827E 02		1.368646E 02				
3	-5.341594E 02		-6.231853E 02		1.513450E 03				
4	-1.905912E 02		-2.223549E 02		5.400093E 02				
5	-9.995206E 01		1.166093E 02		2.831985E 02				
COMPUTED EIGENVALUE AND EIGENVECTOR									
RESIDUAL									
1	7.647050E-01		9.411714E-01		7.46E-07				
	-5.370591E-01		-2.236047E-02						
	-3.185631E-01		-1.325891E-02						
	-3.713154E-01		-1.545933E-02						
	-5.810177E-01		-2.418823E-02						
	-9.991344E-01		-4.160066E-02						
2	-1.000000E 00		-1.333308E 00		4.84E-07				
	-2.592583E-01		-8.606803E-02						
	5.948840E-01		1.974900E-01						
	8.487330E-01		2.817618E-01						
	9.490676E-01		3.150727E-01						
	2.346767E-01		7.791054E-02						
3	-3.529415E-01		-4.117646E-01		6.39E-07				
	2.828103E-02		2.718275E-01						
	-4.315349E-01		-4.140894E-01						
	1.024388E-02		3.613692E-01						
	3.964370E-01		9.180620E-01						
	-2.121114E-01		-5.137083E-01						
4	-3.529405E-01		-4.117613E-01		4.82E-07				
	5.404448E-01		5.624916E-02						
	9.812239E-01		1.928717E-01						
	8.455181E-01		9.440637E-02						
	4.273624E-01		-2.578912E-02						
	-2.842883E-01		7.476088E-03						
5	-3.529399E-01		4.117583E-01		7.78E-07				
	1.844878E-03		9.193128E-02						
	1.966345E-02		9.801325E-01						
	1.144641E-02		5.705664E-01						
	2.005622E-02		9.997990E-01						
	3.575820E-03		1.782405E-01						

Fig. 1

and, if IERR is zero and eigenvectors are desired:

CALL CQZVEC (NM, N, AR, AI, BR, BI, ALFR, ALFI, BETA, ZR, ZI)

where parameter usage is explained in the algorithm that follows.

In timing the subroutines on the IBM 370/195, it was found that they run between two and three times slower than the corresponding real subroutines for systems of the same order. This is felt to be a quite favorable result and is attributable chiefly to two factors: (1) double shifts used with real systems to avoid complex arithmetic can be replaced by more efficient single shifts here, and (2) the amount of complex arithmetic itself can be reduced by choosing one of the defining elements of each 2×2 unitary transformation to be real.

The subroutines provide an alternative to the LZ algorithm developed by Kaufman [2] for the solution of the complex generalized eigenproblem. The elementary transformations employed in the LZ algorithm, even though stabilized, do not preclude possible numerical growth with consequent loss of accuracy; the unitary transformations employed in the QZ algorithm, on the other hand,

are stable. Although unitary transformations require more arithmetic than elementary transformations, speed comparisons between the algorithms are complicated by the fact that the LZ algorithm employs complex operands while the QZ algorithm employs programmed complex arithmetic on real operands. With certain CDC compilers, for example, it has been demonstrated that formal complex arithmetic is more efficient than programmed complex arithmetic, thereby strengthening the advantage of the LZ algorithm; in contrast, with the H compiler on the IBM 370/195, the reverse is true: the complex QZ algorithm actually runs faster than does the LZ algorithm.

A long precision IBM version of the subroutines was first written and tested on the 370/195 at Argonne National Laboratory; the standard single precision Fortran version was then derived from it and also tested on the 370/195. There are no machine-dependent constants in the subroutines, so the standard version should run directly on different machines.

The example chosen for illustration is the same as that used in [2] and appears here as Figure 1. The quantities labeled RESIDUAL are computed as

$$\| \beta_i A z_i - \alpha_i B z_i \|_1 / (\| z_i \|_1 * [\beta_i \| A \|_1 + \alpha_i \| B \|_1])$$

where $\beta_i = \text{BETA}(I)$, $\alpha_i = \text{CMLX}(\text{ALFR}(I), \text{ALFI}(I))$, and $z_i = \text{CMLX}(\text{ZR}(*, I), \text{ZI}(*, I))$.

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2. KAUFMAN, L.C. Algorithm 496. The LZ algorithm to solve the generalized eigenvalue problem for complex matrices. *ACM Trans. Math. Software* 1, 3 (Sept. 1975), 271-281.
3. MOLER, C.B., AND STEWART, G.W. An algorithm for generalized matrix eigenvalue problems. *SIAM J. Numer. Anal.* 10 (April 1973), 241-256.
4. WARD, R.C. The combination shift QZ algorithm. *SIAM J. Numer. Anal.* 12 (Dec. 1975), 835-853.

ALGORITHM

NAME(*n*): indicates a Fortran module with *n* records
 NAME^T: indicates "NAME" is included for testing purposes
 NAME^D: indicates "NAME" contains test data
 Contents: CQZHES(281), CQZVAL(382), CQZVEC(138), TEST^T(75),
 RMATIN^T(65), CGGWZR^T(120), DATA1^D(37), DATA2^D(36)

```

C                                                     HES 10
C -----HES 20
C                                                     HES 30
C   SUBROUTINE CQZHES(NM,N,AR,AI,BR,BI,MATZ,ZR,ZI)      HES 40
C                                                     HES 50
C   INTEGER I,J,K,L,N,K1,LB,L1,NM,NK1,NM1             HES 60
C   REAL AR(NM,N),AI(NM,N),BR(NM,N),BI(NM,N),ZR(NM,N),ZI(NM,N) HES 70
C   REAL R,S,T,TT,U1,U2,XI,XR,YI,YR,RHO,U1I           HES 80
C   REAL SQRT,CABS,ABS                                 HES 90
C   LOGICAL MATZ                                       HES 100
C   COMPLEX CMLX                                       HES 110
C                                                     HES 120
C   THIS SUBROUTINE IS A COMPLEX ANALOGUE OF THE FIRST STEP OF THE HES 130
C   QZ ALGORITHM FOR SOLVING GENERALIZED MATRIX EIGENVALUE PROBLEMS, HES 140
C   SIAM J. NUMER. ANAL. 10, 241-256(1973) BY MOLER AND STEWART. HES 150
C                                                     HES 160
C   THIS SUBROUTINE ACCEPTS A PAIR OF COMPLEX GENERAL MATRICES AND HES 170
C   REDUCES ONE OF THEM TO UPPER HESSENBERG FORM WITH REAL (AND NON- HES 180
C   NEGATIVE) SUBDIAGONAL ELEMENTS AND THE OTHER TO UPPER TRIANGULAR HES 190
C   FORM USING UNITARY TRANSFORMATIONS. IT IS USUALLY FOLLOWED BY HES 200
C   CQZVAL AND POSSIBLY CQZVEC.                       HES 210
C                                                     HES 220
C   ON INPUT-                                          HES 230

```

```

C                                     HES 240
C   NM MUST BE SET TO THE ROW DIMENSION OF TWO-DIMENSIONAL   HES 250
C   ARRAY PARAMETERS AS DECLARED IN THE CALLING PROGRAM       HES 260
C   DIMENSION STATEMENT,                                       HES 270
C                                                             HES 280
C   N IS THE ORDER OF THE MATRICES,                             HES 290
C                                                             HES 300
C   A=(AR,AI) CONTAINS A COMPLEX GENERAL MATRIX,               HES 310
C                                                             HES 320
C   B=(BR,BI) CONTAINS A COMPLEX GENERAL MATRIX,               HES 330
C                                                             HES 340
C   MATZ SHOULD BE SET TO .TRUE. IF THE RIGHT HAND TRANSFORMATIONS HES 350
C   ARE TO BE ACCUMULATED FOR LATER USE IN COMPUTING           HES 360
C   EIGENVECTORS, AND TO .FALSE. OTHERWISE.                   HES 370
C                                                             HES 380
C   ON OUTPUT-                                                 HES 390
C                                                             HES 400
C   A HAS BEEN REDUCED TO UPPER HESSENBERG FORM. THE ELEMENTS HES 410
C   BELOW THE FIRST SUBDIAGONAL HAVE BEEN SET TO ZERO, AND THE HES 420
C   SUBDIAGONAL ELEMENTS HAVE BEEN MADE REAL (AND NON-NEGATIVE), HES 430
C                                                             HES 440
C   B HAS BEEN REDUCED TO UPPER TRIANGULAR FORM. THE ELEMENTS HES 450
C   BELOW THE MAIN DIAGONAL HAVE BEEN SET TO ZERO,             HES 460
C                                                             HES 470
C   Z=(ZR,ZI) CONTAINS THE PRODUCT OF THE RIGHT HAND           HES 480
C   TRANSFORMATIONS IF MATZ HAS BEEN SET TO .TRUE.            HES 490
C   OTHERWISE, Z IS NOT REFERENCED.                            HES 500
C                                                             HES 510
C   QUESTIONS AND COMMENTS SHOULD BE DIRECTED TO B. S. GARBOW, HES 520
C   APPLIED MATHEMATICS DIVISION, ARGONNE NATIONAL LABORATORY HES 530
C   -----HES 550
C   ***** INITIALIZE Z ***** HES 560
C   IF (.NOT. MATZ) GO TO 10 HES 570
C   DO 3 I = 1, N HES 590
C   DO 2 J = 1, N HES 610
C   ZR(I,J) = 0.0 HES 620
C   ZI(I,J) = 0.0 HES 630
C   2 CONTINUE HES 640
C   ZR(I,I) = 1.0 HES 650
C   3 CONTINUE HES 660
C   ***** REDUCE B TO UPPER TRIANGULAR FORM WITH HES 670
C   TEMPORARILY REAL DIAGONAL ELEMENTS ***** HES 680
C   10 IF (N .LE. 1) GO TO 170 HES 690
C   NM1 = N - 1 HES 700
C   DO 100 L = 1, NM1 HES 710
C   L1 = L + 1 HES 720
C   S = 0.0 HES 730
C   DO 20 I = L, N HES 740
C   S = S + ABS(BR(I,L)) + ABS(BI(I,L)) HES 750
C   20 CONTINUE HES 760
C   IF (S .EQ. 0.0) GO TO 100 HES 770
C   RHO = 0.0 HES 780
C   DO 25 I = L, N HES 790
C   BR(I,L) = BR(I,L) / S HES 800
C   BI(I,L) = BI(I,L) / S HES 810
C   RHO = RHO + BR(I,L)**2 + BI(I,L)**2 HES 820
C   25 CONTINUE HES 830
C   R = Sqrt(RHO) HES 840
C   XR = CABS(CMPLX(BR(L,L),BI(L,L))) HES 850
C   IF (XR .EQ. 0.0) GO TO 27 HES 860
C   RHO = RHO + XR * R HES 870
C   U1 = -BR(L,L) / XR HES 880
C   U11 = -BI(L,L) / XR HES 890
C   YR = R / XR + 1.0 HES 900
C   BR(L,L) = YR * BR(L,L) HES 910
C   BI(L,L) = YR * BI(L,L) HES 920

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```

      GO TO 28
C
27  BR(L,L) = R
    U1 = -1.0
    U1I = 0.0
C
28  DO 50 J = L1, N
    T = 0.0
    TI = 0.0
C
    DO 30 I = L, N
      T = T + BR(I,L) * BR(I,J) + BI(I,L) * BI(I,J)
      TI = TI + BR(I,L) * BI(I,J) - BI(I,L) * BR(I,J)
30  CONTINUE
C
    T = T / RHO
    TI = TI / RHO
C
    DO 40 I = L, N
      BR(I,J) = BR(I,J) - T * BR(I,L) + TI * BI(I,L)
      BI(I,J) = BI(I,J) - T * BI(I,L) - TI * BR(I,L)
40  CONTINUE
C
    XI = U1 * BI(L,J) - U1I * BR(L,J)
    BR(L,J) = U1 * BR(L,J) + U1I * BI(L,J)
    BI(L,J) = XI
50  CONTINUE
C
    DO 80 J = 1, N
    T = 0.0
    TI = 0.0
C
    DO 60 I = L, N
      T = T + BR(I,L) * AR(I,J) + BI(I,L) * AI(I,J)
      TI = TI + BR(I,L) * AI(I,J) - BI(I,L) * AR(I,J)
60  CONTINUE
C
    T = T / RHO
    TI = TI / RHO
C
    DO 70 I = L, N
      AR(I,J) = AR(I,J) - T * BR(I,L) + TI * BI(I,L)
      AI(I,J) = AI(I,J) - T * BI(I,L) - TI * BR(I,L)
70  CONTINUE
C
    XI = U1 * AI(L,J) - U1I * AR(L,J)
    AR(L,J) = U1 * AR(L,J) + U1I * AI(L,J)
    AI(L,J) = XI
80  CONTINUE
C
    BR(L,L) = R * S
    BI(L,L) = 0.0
C
    DO 90 I = L1, N
      BR(I,L) = 0.0
      BI(I,L) = 0.0
90  CONTINUE
C
100 CONTINUE
C ***** REDUCE A TO UPPER HESSENBERG FORM WITH REAL SUBDIAGONAL
C ***** ELEMENTS, WHILE KEEPING B TRIANGULAR *****
DO 160 K = 1, NM1
  K1 = K + 1
C ***** SET BOTTOM ELEMENT IN K-TH COLUMN OF A REAL *****
  IF (AI(N,K) .EQ. 0.0) GO TO 105
  R = CABS(CPLX(AR(N,K),AI(N,K)))
  U1 = AR(N,K) / R
  U1I = AI(N,K) / R
  AR(N,K) = R
  AI(N,K) = 0.0
C
  DO 103 J = K1, N
    XI = U1 * AI(N,J) - U1I * AR(N,J)
    AR(N,J) = U1 * AR(N,J) + U1I * AI(N,J)
    AI(N,J) = XI
103 CONTINUE

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HES 1000
 HES 1010
 HES 1020
 HES 1030
 HES 1040
 HES 1050
 HES 1060
 HES 1070
 HES 1080
 HES 1090
 HES 1100
 HES 1110
 HES 1120
 HES 1130
 HES 1140
 HES 1150
 HES 1160
 HES 1170
 HES 1180
 HES 1190
 HES 1200
 HES 1210
 HES 1220
 HES 1230
 HES 1240
 HES 1250
 HES 1260
 HES 1270
 HES 1280
 HES 1290
 HES 1300
 HES 1310
 HES 1320
 HES 1330
 HES 1340
 HES 1350
 HES 1360
 HES 1370
 HES 1380
 HES 1390
 HES 1400
 HES 1410
 HES 1420
 HES 1430
 HES 1440
 HES 1450
 HES 1460
 HES 1470
 HES 1480
 HES 1490
 HES 1500
 HES 1510
 HES 1520
 HES 1530
 HES 1540
 HES 1550
 HES 1560
 HES 1570
 HES 1580
 HES 1590
 HES 1600
 HES 1610
 HES 1620
 HES 1630
 HES 1640
 HES 1650
 HES 1660
 HES 1670
 HES 1680
 HES 1690
 HES 1700
 HES 1710
 HES 1720
 HES 1730
 HES 1740
 HES 1750

C		HES 1760
	XI = U1 * BI(N,N) - U1I * BR(N,N)	HES 1770
	BR(N,N) = U1 * BR(N,N) + U1I * BI(N,N)	HES 1780
	BI(N,N) = XI	HES 1790
105	IF (K .EQ. NM1) GO TO 170	HES 1800
	NK1 = NM1 - K	HES 1810
C	***** FOR L=N-1 STEP -1 UNTIL K+1 DO -- *****	HES 1820
	DO 150 LB = 1, NK1	HES 1830
	L = N - LB	HES 1840
	L1 = L + 1	HES 1850
C	***** ZERO A(L+1,K) *****	HES 1860
	S = ABS(AR(L,K)) + ABS(AI(L,K)) + AR(L1,K)	HES 1870
	IF (S .EQ. 0.0) GO TO 150	HES 1880
	U1 = AR(L,K) / S	HES 1890
	U1I = AI(L,K) / S	HES 1900
	U2 = AR(L1,K) / S	HES 1910
	R = SQRT(U1*U1+U1I*U1I+U2*U2)	HES 1920
	U1 = U1 / R	HES 1930
	U1I = U1I / R	HES 1940
	U2 = U2 / R	HES 1950
	AR(L,K) = R * S	HES 1960
	AI(L,K) = 0.0	HES 1970
	AR(L1,K) = 0.0	HES 1980
C		HES 1990
	DO 110 J = K1, N	HES 2000
	XR = AR(L,J)	HES 2010
	XI = AI(L,J)	HES 2020
	YR = AR(L1,J)	HES 2030
	YI = AI(L1,J)	HES 2040
	AR(L,J) = U1 * XR + U1I * XI + U2 * YR	HES 2050
	AI(L,J) = U1 * XI - U1I * XR + U2 * YI	HES 2060
	AR(L1,J) = U1 * YR - U1I * YI - U2 * XR	HES 2070
	AI(L1,J) = U1 * YI + U1I * YR - U2 * XI	HES 2080
110	CONTINUE	HES 2090
C		HES 2100
	XR = BR(L,L)	HES 2110
	BR(L,L) = U1 * XR	HES 2120
	BI(L,L) = -U1I * XR	HES 2130
	BR(L1,L) = -U2 * XR	HES 2140
C		HES 2150
	DO 120 J = L1, N	HES 2160
	XR = BR(L,J)	HES 2170
	XI = BI(L,J)	HES 2180
	YR = BR(L1,J)	HES 2190
	YI = BI(L1,J)	HES 2200
	BR(L,J) = U1 * XR + U1I * XI + U2 * YR	HES 2210
	BI(L,J) = U1 * XI - U1I * XR + U2 * YI	HES 2220
	BR(L1,J) = U1 * YR - U1I * YI - U2 * XR	HES 2230
	BI(L1,J) = U1 * YI + U1I * YR - U2 * XI	HES 2240
120	CONTINUE	HES 2250
C	***** ZERO B(L+1,L) *****	HES 2260
	S = ABS(BR(L1,L1)) + ABS(BI(L1,L1)) + ABS(BR(L1,L))	HES 2270
	IF (S .EQ. 0.0) GO TO 150	HES 2280
	U1 = BR(L1,L1) / S	HES 2290
	U1I = BI(L1,L1) / S	HES 2300
	U2 = BR(L1,L) / S	HES 2310
	R = SQRT(U1*U1+U1I*U1I+U2*U2)	HES 2320
	U1 = U1 / R	HES 2330
	U1I = U1I / R	HES 2340
	U2 = U2 / R	HES 2350
	BR(L1,L1) = R * S	HES 2360
	BI(L1,L1) = 0.0	HES 2370
	BR(L1,L) = 0.0	HES 2380
C		HES 2390
	DO 130 I = 1, L	HES 2400
	XR = BR(I,L1)	HES 2410
	XI = BI(I,L1)	HES 2420
	YR = BR(I,L)	HES 2430
	YI = BI(I,L)	HES 2440
	BR(I,L1) = U1 * XR + U1I * XI + U2 * YR	HES 2450
	BI(I,L1) = U1 * XI - U1I * XR + U2 * YI	HES 2460
	BR(I,L) = U1 * YR - U1I * YI - U2 * XR	HES 2470
	BI(I,L) = U1 * YI + U1I * YR - U2 * XI	HES 2480
130	CONTINUE	HES 2490
C		HES 2500
	DO 140 I = 1, N	HES 2510

	XR = AR(I,L1)	HES 2520
	XI = AI(I,L1)	HES 2530
	YR = AR(I,L)	HES 2540
	YI = AI(I,L)	HES 2550
	AR(I,L1) = U1 * XR + U1I * XI + U2 * YR	HES 2560
	AI(I,L1) = U1 * XI - U1I * XR + U2 * YI	HES 2570
	AR(I,L) = U1 * YR - U1I * YI - U2 * XR	HES 2580
	AI(I,L) = U1 * YI + U1I * YR - U2 * XI	HES 2590
140	CONTINUE	HES 2600
C		HES 2610
	IF (.NOT. MATZ) GO TO 150	HES 2620
C		HES 2630
	DO 145 I = 1, N	HES 2640
	XR = ZR(I,L1)	HES 2650
	XI = ZI(I,L1)	HES 2660
	YR = ZR(I,L)	HES 2670
	YI = ZI(I,L)	HES 2680
	ZR(I,L1) = U1 * XR + U1I * XI + U2 * YR	HES 2690
	ZI(I,L1) = U1 * XI - U1I * XR + U2 * YI	HES 2700
	ZR(I,L) = U1 * YR - U1I * YI - U2 * XR	HES 2710
	ZI(I,L) = U1 * YI + U1I * YR - U2 * XI	HES 2720
145	CONTINUE	HES 2730
C		HES 2740
150	CONTINUE	HES 2750
C		HES 2760
160	CONTINUE	HES 2770
C		HES 2780
170	RETURN	HES 2790
C	***** LAST CARD OF CQZHE *****	HES 2800
	END	HES 2810
C		VAL 10
C	-----	VAL 20
C		VAL 30
	SUBROUTINE CQZVAL(NM,N,AR,AI,BR,BI,EPS1,ALFR,ALFI,BETA,	VAL 40
X	MATZ,ZR,ZI,IERR)	VAL 50
C		VAL 60
	INTEGER I,J,K,L,N,EN,K1,K2,LL,L1,NA,NM,ITS,KM1,LM1,	VAL 70
X	ENM2,IERR,LOR1,ENORN	VAL 80
	REAL AR(NM,N),AI(NM,N),BR(NM,N),BI(NM,N),ALFR(N),ALFI(N),	VAL 90
X	BETA(N),ZR(NM,N),ZI(NM,N)	VAL 100
	REAL R,S,A1,A2,EP,SH,U1,U2,XI,XR,YI,ANI,A1I,A33,A34,A43,A44,	VAL 110
X	BNI,B11,B33,B44,SHI,U1I,A33I,A34I,A43I,A44I,B33I,B44I,	VAL 120
X	EPSA,EP SB,EPS1,ANORM,BNORM,B3344,B3344I	VAL 130
	REAL SQRT,CABS,ABS	VAL 140
	INTEGER MAXO	VAL 150
	LOGICAL MATZ	VAL 160
	COMPLEX Z3	VAL 170
	COMPLEX CSQRT,CMPLX	VAL 180
	REAL REAL,AIMAG	VAL 190
C		VAL 200
C		VAL 210
C		VAL 220
C		VAL 230
C		VAL 240
C	THIS SUBROUTINE IS A COMPLEX ANALOGUE OF STEPS 2 AND 3 OF THE	VAL 250
C	QZ ALGORITHM FOR SOLVING GENERALIZED MATRIX EIGENVALUE PROBLEMS,	VAL 260
C	SIAM J. NUMER. ANAL. 10, 241-256(1973) BY MOLER AND STEWART,	VAL 270
C	AS MODIFIED IN TECHNICAL NOTE NASA TN E-7305(1973) BY WARD.	VAL 280
C		VAL 290
C	THIS SUBROUTINE ACCEPTS A PAIR OF COMPLEX MATRICES, ONE OF THEM	VAL 300
C	IN UPPER HESSENBERG FORM AND THE OTHER IN UPPER TRIANGULAR FORM,	VAL 310
C	THE HESSENBERG MATRIX MUST FURTHER HAVE REAL SUBDIAGONAL ELEMENTS.	VAL 320
C	IT REDUCES THE HESSENBERG MATRIX TO TRIANGULAR FORM USING	VAL 330
C	UNITARY TRANSFORMATIONS WHILE MAINTAINING THE TRIANGULAR FORM	VAL 340
C	OF THE OTHER MATRIX AND FURTHER MAKING ITS DIAGONAL ELEMENTS	VAL 350
C	REAL AND NON-NEGATIVE. IT THEN RETURNS QUANTITIES WHOSE RATIOS	VAL 360
C	GIVE THE GENERALIZED EIGENVALUES. IT IS USUALLY PRECEDED BY	VAL 370
C	CQZHE AND POSSIBLY FOLLOWED BY CQZVEC.	VAL 380
C		VAL 390
C	ON INPUT-	VAL 400
C		VAL 410
C	NM MUST BE SET TO THE ROW DIMENSION OF TWO-DIMENSIONAL	VAL 420
C	ARRAY PARAMETERS AS DECLARED IN THE CALLING PROGRAM	VAL 430

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C          DIMENSION STATEMENT, VAL 440
C          VAL 450
C          N IS THE ORDER OF THE MATRICES, VAL 460
C          VAL 470
C          A=(AR,AI) CONTAINS A COMPLEX UPPER HESSENBERG MATRIX VAL 480
C          WITH REAL SUBDIAGONAL ELEMENTS, VAL 490
C          VAL 500
C          B=(BR,BI) CONTAINS A COMPLEX UPPER TRIANGULAR MATRIX, VAL 510
C          VAL 520
C          EPS1 IS A TOLERANCE USED TO DETERMINE NEGLIGIBLE ELEMENTS. VAL 530
C          EPS1 = 0.0 (OR NEGATIVE) MAY BE INPUT, IN WHICH CASE AN VAL 540
C          ELEMENT WILL BE NEGLECTED ONLY IF IT IS LESS THAN ROUND OFF VAL 550
C          ERROR TIMES THE NORM OF ITS MATRIX. IF THE INPUT EPS1 IS VAL 560
C          POSITIVE, THEN AN ELEMENT WILL BE CONSIDERED NEGLIGIBLE VAL 570
C          IF IT IS LESS THAN EPS1 TIMES THE NORM OF ITS MATRIX. A VAL 580
C          POSITIVE VALUE OF EPS1 MAY RESULT IN FASTER EXECUTION, VAL 590
C          BUT LESS ACCURATE RESULTS, VAL 600
C          VAL 610
C          MATZ SHOULD BE SET TO .TRUE. IF THE RIGHT HAND TRANSFORMATIONS VAL 620
C          ARE TO BE ACCUMULATED FOR LATER USE IN COMPUTING VAL 630
C          EIGENVECTORS, AND TO .FALSE. OTHERWISE, VAL 640
C          VAL 650
C          Z=(ZR,ZI) CONTAINS, IF MATZ HAS BEEN SET TO .TRUE., THE VAL 660
C          TRANSFORMATION MATRIX PRODUCED IN THE REDUCTION VAL 670
C          BY CQZHES, IF PERFORMED, OR ELSE THE IDENTITY MATRIX. VAL 680
C          IF MATZ HAS BEEN SET TO .FALSE., Z IS NOT REFERENCED. VAL 690
C          VAL 700
C          ON OUTPUT- VAL 710
C          VAL 720
C          A HAS BEEN REDUCED TO UPPER TRIANGULAR FORM. THE ELEMENTS VAL 730
C          BELOW THE MAIN DIAGONAL HAVE BEEN SET TO ZERO, VAL 740
C          VAL 750
C          B IS STILL IN UPPER TRIANGULAR FORM, ALTHOUGH ITS ELEMENTS VAL 760
C          HAVE BEEN ALTERED. IN PARTICULAR, ITS DIAGONAL HAS BEEN SET VAL 770
C          REAL AND NON-NEGATIVE. THE LOCATION BR(N,1) IS USED TO VAL 780
C          STORE EPS1 TIMES THE NORM OF B FOR LATER USE BY CQZVEC, VAL 790
C          VAL 800
C          ALFR AND ALFI CONTAIN THE REAL AND IMAGINARY PARTS OF THE VAL 810
C          DIAGONAL ELEMENTS OF THE TRIANGULARIZED A MATRIX, VAL 820
C          VAL 830
C          BETA CONTAINS THE REAL NON-NEGATIVE DIAGONAL ELEMENTS OF THE VAL 840
C          CORRESPONDING B. THE GENERALIZED EIGENVALUES ARE THEN VAL 850
C          THE RATIOS ((ALFR+I*ALFI)/BETA), VAL 860
C          VAL 870
C          Z CONTAINS THE PRODUCT OF THE RIGHT HAND TRANSFORMATIONS VAL 880
C          (FOR BOTH STEPS) IF MATZ HAS BEEN SET TO .TRUE., VAL 890
C          VAL 900
C          IERR IS SET TO VAL 910
C          ZERO FOR NORMAL RETURN, VAL 920
C          J IF AR(J,J-1) HAS NOT BECOME VAL 930
C          ZERO AFTER 50 ITERATIONS. VAL 940
C          VAL 950
C          QUESTIONS AND COMMENTS SHOULD BE DIRECTED TO B. S. GARBOW, VAL 960
C          APPLIED MATHEMATICS DIVISION, ARGONNE NATIONAL LABORATORY VAL 970
C          VAL 980
C          ----- VAL 990
C          VAL 1000
C          IERR = 0 VAL 1010
C          ***** COMPUTE EPSA,EPSB ***** VAL 1020
C          ANORM = 0.0 VAL 1030
C          BNORM = 0.0 VAL 1040
C          VAL 1050
C          DO 30 I = 1, N VAL 1060
C          ANI = 0.0 VAL 1070
C          IF (I .NE. 1) ANI = ABS(AR(I,I-1)) VAL 1080
C          BNI = 0.0 VAL 1090
C          VAL 1100
C          DO 20 J = 1, N VAL 1110
C          ANI = ANI + ABS(AR(I,J)) + ABS(AI(I,J)) VAL 1120
C          BNI = BNI + ABS(BR(I,J)) + ABS(BI(I,J)) VAL 1130
C          20 CONTINUE VAL 1140
C          VAL 1150
C          IF (ANI .GT. ANORM) ANORM = ANI VAL 1160
C          IF (BNI .GT. BNORM) BNORM = BNI VAL 1170
C          30 CONTINUE VAL 1180
C          VAL 1190

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      IF (ANORM .EQ. 0.0) ANORM = 1.0
      IF (BNORM .EQ. 0.0) BNORM = 1.0
      EP = EPS1
      IF (EP .GT. 0.0) GO TO 50
C     ***** COMPUTE ROUND OFF LEVEL IF EPS1 IS ZERO *****
      EP = 1.0
40    EP = EP / 2.0
      IF (1.0 + EP .GT. 1.0) GO TO 40
50    EPSA = EP * ANORM
      EPSB = EP * BNORM
C     ***** REDUCE A TO TRIANGULAR FORM, WHILE
C     KEEPING B TRIANGULAR *****
      LOR1 = 1
      ENORN = N
      EN = N
C     ***** BEGIN QZ STEP *****
60    IF (EN .EQ. 0) GO TO 1001
      IF (.NOT. MATZ) ENORN = EN
      ITS = 0
      NA = EN - 1
      ENM2 = NA - 1
C     ***** CHECK FOR CONVERGENCE OR REDUCIBILITY.
C     FOR L=EN STEP -1 UNTIL 1 DO -- *****
70    DO 80 LL = 1, EN
      LM1 = EN - LL
      L = LM1 + 1
      IF (L .EQ. 1) GO TO 95
      IF (ABS(AR(L,LM1)) .LE. EPSA) GO TO 90
80    CONTINUE
C
90    AR(L,LM1) = 0.0
C     ***** SET DIAGONAL ELEMENT AT TOP OF B REAL *****
95    B11 = CABS(CMPLX(BR(L,L),BI(L,L)))
      IF (BI(L,L) .EQ. 0.0) GO TO 98
      U1 = BR(L,L) / B11
      U1I = BI(L,L) / B11
C
      DO 97 J = L, ENORN
      XI = U1 * AI(L,J) - U1I * AR(L,J)
      AR(L,J) = U1 * AR(L,J) + U1I * AI(L,J)
      AI(L,J) = XI
      XI = U1 * BI(L,J) - U1I * BR(L,J)
      BR(L,J) = U1 * BR(L,J) + U1I * BI(L,J)
      BI(L,J) = XI
97    CONTINUE
C
      BI(L,L) = 0.0
98    IF (L .NE. EN) GO TO 100
C     ***** 1-BY-1 BLOCK ISOLATED *****
      ALFR(EN) = AR(EN,EN)
      ALFI(EN) = AI(EN,EN)
      BETA(EN) = B11
      EN = NA
      GO TO 60
C     ***** CHECK FOR SMALL TOP OF B *****
100   L1 = L + 1
      IF (B11 .GT. EPSB) GO TO 120
      BR(L,L) = 0.0
      S = ABS(AR(L,L)) + ABS(AI(L,L)) + ABS(AR(L1,L))
      U1 = AR(L,L) / S
      U1I = AI(L,L) / S
      U2 = AR(L1,L) / S
      R = SQRT(U1*U1+U1I*U1I+U2*U2)
      U1 = U1 / R
      U1I = U1I / R
      U2 = U2 / R
      AR(L,L) = R * S
      AI(L,L) = 0.0
C
      DO 110 J = L1, ENORN
      XR = AR(L,J)
      XI = AI(L,J)
      YR = AR(L1,J)
      YI = AI(L1,J)
      AR(L,J) = U1 * XR + U1I * XI + U2 * YR
      AI(L,J) = U1 * XI - U1I * XR + U2 * YI

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VAL 1200
 VAL 1210
 VAL 1220
 VAL 1230
 VAL 1240
 VAL 1250
 VAL 1260
 VAL 1270
 VAL 1280
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 VAL 1300
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 VAL 1370
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 VAL 1760
 VAL 1770
 VAL 1780
 VAL 1790
 VAL 1800
 VAL 1810
 VAL 1820
 VAL 1830
 VAL 1840
 VAL 1850
 VAL 1860
 VAL 1870
 VAL 1880
 VAL 1890
 VAL 1900
 VAL 1910
 VAL 1920
 VAL 1930
 VAL 1940
 VAL 1950

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      AR(L1,J) = U1 * YR - U1I * YI - U2 * XR          VAL 1960
      AI(L1,J) = U1 * YI + U1I * YR - U2 * XI          VAL 1970
      XR = BR(L,J)                                     VAL 1980
      XI = BI(L,J)                                     VAL 1990
      YR = BR(L1,J)                                    VAL 2000
      YI = BI(L1,J)                                    VAL 2010
      BR(L1,J) = U1 * YR - U1I * YI - U2 * XR          VAL 2040
      BR(L,J) = U1 * XR + U1I * XI + U2 * YR           VAL 2020
      BI(L,J) = U1 * XI - U1I * XR + U2 * YI           VAL 2030
      BI(L1,J) = U1 * YI + U1I * YR - U2 * XI          VAL 2050
110 CONTINUE                                          VAL 2060
C                                                      VAL 2070
      LM1 = L                                          VAL 2080
      L = L1                                           VAL 2090
      GO TO 90                                          VAL 2100
C ***** ITERATION STRATEGY *****                 VAL 2110
120 IF (ITS .EQ. 50) GO TO 1000                       VAL 2120
      IF (ITS .EQ. 10) GO TO 135                      VAL 2130
C ***** DETERMINE SHIFT *****                   VAL 2140
      B33 = BR(NA,NA)                                  VAL 2150
      B33I = BI(NA,NA)                                 VAL 2160
      IF (CABS(CMPLX(B33,B33I)) .GE. EPSB) GO TO 122    VAL 2170
      B33 = EPSB                                       VAL 2180
      B33I = 0.0                                        VAL 2190
122 B44 = BR(EN,EN)                                    VAL 2200
      B44I = BI(EN,EN)                                  VAL 2210
      IF (CABS(CMPLX(B44,B44I)) .GE. EPSB) GO TO 124    VAL 2220
      B44 = EPSB                                       VAL 2230
      B44I = 0.0                                        VAL 2240
124 B3344 = B33 * B44 - B33I * B44I                  VAL 2250
      B3344I = B33 * B44I + B33I * B44                VAL 2260
      A33 = AR(NA,NA) * B44 - AI(NA,NA) * B44I        VAL 2270
      A33I = AR(NA,NA) * B44I + AI(NA,NA) * B44        VAL 2280
      A34 = AR(NA,EN) * B33 - AI(NA,EN) * B33I        VAL 2290
      X   = AR(NA,NA) * BR(NA,EN) + AI(NA,NA) * BI(NA,EN) VAL 2300
      A34I = AR(NA,EN) * B33I + AI(NA,EN) * B33        VAL 2310
      X   = AR(NA,NA) * BI(NA,EN) - AI(NA,NA) * BR(NA,EN) VAL 2320
      A43 = AR(EN,NA) * B44                             VAL 2330
      A43I = AR(EN,NA) * B44I                           VAL 2340
      A44 = AR(EN,EN) * B33 - AI(EN,EN) * B33I - AR(EN,NA) * BR(NA,EN) VAL 2350
      A44I = AR(EN,EN) * B33I + AI(EN,EN) * B33 - AR(EN,NA) * BI(NA,EN) VAL 2360
      SH = A44                                           VAL 2370
      SHI = A44I                                         VAL 2380
      XR = A34 * A43 - A34I * A43I                     VAL 2390
      XI = A34 * A43I + A34I * A43                     VAL 2400
      IF (XR .EQ. 0.0 .AND. XI .EQ. 0.0) GO TO 140      VAL 2410
      YR = (A33 - SH) / 2.0                             VAL 2420
      YI = (A33I - SHI) / 2.0                           VAL 2430
      Z3 = CSQRT(CMPLX(YR**2-YI**2+XR,2.0*YR*YI+XI))    VAL 2440
      U1 = REAL(Z3)                                       VAL 2450
      U1I = AIMAG(Z3)                                    VAL 2460
      IF (YR * U1 + YI * U1I .GE. 0.0) GO TO 125        VAL 2470
      U1 = -U1                                           VAL 2480
      U1I = -U1I                                         VAL 2490
125 Z3 = (CMPLX(SH,SHI) - CMPLX(XR,XI) / CMPLX(YR+U1,YI+U1I)) VAL 2500
      X   / CMPLX(B3344,B3344I)                          VAL 2510
      SH = REAL(Z3)                                       VAL 2520
      SHI = AIMAG(Z3)                                    VAL 2530
      GO TO 140                                          VAL 2540
C ***** AD HOC SHIFT *****                       VAL 2550
135 SH = AR(EN,NA) + AR(NA,ENM2)                       VAL 2560
      SHI = 0.0                                          VAL 2570
C ***** DETERMINE ZEROth COLUMN OF A *****       VAL 2580
140 A1 = AR(L,L) / B11 - SH                             VAL 2590
      A1I = AI(L,L) / B11 - SHI                         VAL 2600
      A2 = AR(L1,L) / B11                               VAL 2610
      ITS = ITS + 1                                     VAL 2620
      IF (.NOT. MATZ) LOR1 = L                          VAL 2630
C ***** MAIN LOOP *****                         VAL 2640
      DO 260 K = L, NA                                  VAL 2650
          K1 = K + 1                                     VAL 2660
          K2 = K + 2                                     VAL 2670
          KM1 = MAX0(K-1,L)                             VAL 2680
C ***** ZERO A(K+1,K-1) *****                   VAL 2690
          IF (K .EQ. L) GO TO 170                       VAL 2700
          A1 = AR(K,KM1)                                 VAL 2710

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	AI = AI(K,KM1)	VAL 2720
	A2 = AR(K1,KM1)	VAL 2730
170	S = ABS(A1) + ABS(A11) + ABS(A2)	VAL 2740
	U1 = A1 / S	VAL 2750
	U11 = A11 / S	VAL 2760
	U2 = A2 / S	VAL 2770
	R = SQRT(U1*U1+U11*U11+U2*U2)	VAL 2780
	U1 = U1 / R	VAL 2790
	U11 = U11 / R	VAL 2800
	U2 = U2 / R	VAL 2810
C		VAL 2820
	DO 180 J = KM1, ENORN	VAL 2830
	XR = AR(K,J)	VAL 2840
	XI = AI(K,J)	VAL 2850
	YR = AR(K1,J)	VAL 2860
	YI = AI(K1,J)	VAL 2870
	AR(K,J) = U1 * XR + U11 * XI + U2 * YR	VAL 2880
	AI(K,J) = U1 * XI - U11 * XR + U2 * YI	VAL 2890
	AR(K1,J) = U1 * YR - U11 * YI - U2 * XR	VAL 2900
	AI(K1,J) = U1 * YI + U11 * YR - U2 * XI	VAL 2910
	XR = BR(K,J)	VAL 2920
	XI = BI(K,J)	VAL 2930
	YR = BR(K1,J)	VAL 2940
	YI = BI(K1,J)	VAL 2950
	BR(K,J) = U1 * XR + U11 * XI + U2 * YR	VAL 2960
	BI(K,J) = U1 * XI - U11 * XR + U2 * YI	VAL 2970
	BR(K1,J) = U1 * YR - U11 * YI - U2 * XR	VAL 2980
	BI(K1,J) = U1 * YI + U11 * YR - U2 * XI	VAL 2990
180	CONTINUE	VAL 3000
C		VAL 3010
	IF (K .EQ. L) GO TO 240	VAL 3020
	AI(K,KM1) = 0.0	VAL 3030
	AR(K1,KM1) = 0.0	VAL 3040
	AI(K1,KM1) = 0.0	VAL 3050
C	***** ZERO B(K+1,K) *****	VAL 3060
240	S = ABS(BR(K1,K1)) + ABS(BI(K1,K1)) + ABS(BR(K1,K))	VAL 3070
	U1 = BR(K1,K1) / S	VAL 3080
	U11 = BI(K1,K1) / S	VAL 3090
	U2 = BR(K1,K) / S	VAL 3100
	R = SQRT(U1*U1+U11*U11+U2*U2)	VAL 3110
	U1 = U1 / R	VAL 3120
	U11 = U11 / R	VAL 3130
	U2 = U2 / R	VAL 3140
	IF (K .EQ. NA) GO TO 245	VAL 3150
	XR = AR(K2,K1)	VAL 3160
	AR(K2,K1) = U1 * XR	VAL 3170
	AI(K2,K1) = -U11 * XR	VAL 3180
	AR(K2,K) = -U2 * XR	VAL 3190
C		VAL 3200
245	DO 250 I = LOR1, K1	VAL 3210
	XR = AR(I,K1)	VAL 3220
	XI = AI(I,K1)	VAL 3230
	YR = AR(I,K)	VAL 3240
	YI = AI(I,K)	VAL 3250
	AR(I,K1) = U1 * XR + U11 * XI + U2 * YR	VAL 3260
	AI(I,K1) = U1 * XI - U11 * XR + U2 * YI	VAL 3270
	AR(I,K) = U1 * YR - U11 * YI - U2 * XR	VAL 3280
	AI(I,K) = U1 * YI + U11 * YR - U2 * XI	VAL 3290
	XR = BR(I,K1)	VAL 3300
	XI = BI(I,K1)	VAL 3310
	YR = BR(I,K)	VAL 3320
	YI = BI(I,K)	VAL 3330
	BR(I,K1) = U1 * XR + U11 * XI + U2 * YR	VAL 3340
	BI(I,K1) = U1 * XI - U11 * XR + U2 * YI	VAL 3350
	BR(I,K) = U1 * YR - U11 * YI - U2 * XR	VAL 3360
	BI(I,K) = U1 * YI + U11 * YR - U2 * XI	VAL 3370
250	CONTINUE	VAL 3380
C		VAL 3390
	BI(K1,K1) = 0.0	VAL 3400
	BR(K1,K) = 0.0	VAL 3410
	BI(K1,K) = 0.0	VAL 3420
	IF (.NOT. MATZ) GO TO 260	VAL 3430
C		VAL 3440
	DO 255 I = 1, N	VAL 3450
	XR = ZR(I,K1)	VAL 3460
	XI = ZI(I,K1)	VAL 3470

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        YR = ZR(I,K)                                VAL 3480
        YI = ZI(I,K)                                VAL 3490
        ZR(I,K1) = U1 * XR + U1I * XI + U2 * YR      VAL 3500
        ZI(I,K1) = U1 * XI - U1I * XR + U2 * YI      VAL 3510
        ZR(I,K) = U1 * YR - U1I * YI - U2 * XR      VAL 3520
        ZI(I,K) = U1 * YI + U1I * YR - U2 * XI      VAL 3530
255   CONTINUE                                      VAL 3540
C                                           VAL 3550
260 CONTINUE                                      VAL 3560
C ***** SET LAST A SUBDIAGONAL REAL AND END QZ STEP ***** VAL 3570
  IF (AI(EN,NA) .EQ. 0.0) GO TO 70
  R = CABS(CMPLX(AR(EN,NA),AI(EN,NA)))              VAL 3580
  U1 = AR(EN,NA) / R                               VAL 3590
  U1I = AI(EN,NA) / R                              VAL 3600
  AR(EN,NA) = R                                    VAL 3610
  AI(EN,NA) = 0.0                                  VAL 3620
C                                           VAL 3630
  DO 270 J = EN, ENORN                             VAL 3640
    XI = U1 * AI(EN,J) - U1I * AR(EN,J)            VAL 3650
    AR(EN,J) = U1 * AR(EN,J) + U1I * AI(EN,J)      VAL 3660
    AI(EN,J) = XI                                  VAL 3670
    XI = U1 * BI(EN,J) - U1I * BR(EN,J)            VAL 3680
    BR(EN,J) = U1 * BR(EN,J) + U1I * BI(EN,J)      VAL 3690
    BI(EN,J) = XI                                  VAL 3700
  270 CONTINUE                                      VAL 3710
C                                           VAL 3720
  GO TO 70                                          VAL 3730
C ***** SET ERROR -- BOTTOM SUBDIAGONAL ELEMENT HAS NOT VAL 3740
C ***** BECOME NEGLIGIBLE AFTER 50 ITERATIONS ***** VAL 3750
1000 IERR = EN                                     VAL 3760
C ***** SAVE EPSB FOR USE BY CQZVEC ***** VAL 3770
1001 IF (N .GT. 1) BR(N,1) = EPSB                 VAL 3780
      RETURN                                       VAL 3790
C ***** LAST CARD OF CQZVAL ***** VAL 3800
      END                                         VAL 3810
                                           VAL 3820

C                                           VEC 10
C ----- VEC 20
C SUBROUTINE CQZVEC(NM,N,AR,AI,BR,BI,ALFR,ALFI,BETA,ZR,ZI) VEC 30
C                                           VEC 40
C INTEGER I,J,K;M,N,EN,II,JJ,NA,NM,NN            VEC 50
C REAL AR(NM,N),AI(NM,N),BR(NM,N),BI(NM,N),ALFR(N),ALFI(N), VEC 60
C X BETA(N),ZR(NM,N),ZI(NM,N)                   VEC 70
C REAL R,T,RI,TI,XI,ALMI,ALMR,BETM,EPSB         VEC 80
C REAL CABS                                     VEC 90
C COMPLEX Z3                                    VEC 100
C COMPLEX CMPLX                                VEC 110
C REAL REAL,AIMAG                              VEC 120
C                                           VEC 130
C                                           VEC 140
C                                           VEC 150
C                                           VEC 160
C                                           VEC 170
C                                           VEC 180
C THIS SUBROUTINE IS A COMPLEX ANALOGUE OF THE FOURTH STEP OF THE VEC 190
C QZ ALGORITHM FOR SOLVING GENERALIZED MATRIX EIGENVALUE PROBLEMS, VEC 200
C SIAM J. NUMER. ANAL. 10, 241-256(1973) BY MOLER AND STEWART. VEC 210
C                                           VEC 220
C THIS SUBROUTINE ACCEPTS A PAIR OF COMPLEX MATRICES IN UPPER VEC 230
C TRIANGULAR FORM, WHERE ONE OF THEM FURTHER MUST HAVE REAL DIAGONAL VEC 240
C ELEMENTS. IT COMPUTES THE EIGENVECTORS OF THE TRIANGULAR PROBLEM VEC 250
C AND TRANSFORMS THE RESULTS BACK TO THE ORIGINAL COORDINATE SYSTEM. VEC 260
C IT IS USUALLY PRECEDED BY CQZHS AND CQZVAL. VEC 270
C                                           VEC 280
C ON INPUT- VEC 290
C                                           VEC 300
C NM MUST BE SET TO THE ROW DIMENSION OF TWO-DIMENSIONAL VEC 310
C ARRAY PARAMETERS AS DECLARED IN THE CALLING PROGRAM VEC 320
C DIMENSION STATEMENT, VEC 330
C                                           VEC 340
C N IS THE ORDER OF THE MATRICES, VEC 350
C                                           VEC 360
C A=(AR,AI) CONTAINS A COMPLEX UPPER TRIANGULAR MATRIX, VEC 370
C                                           VEC 380
C B=(BR,BI) CONTAINS A COMPLEX UPPER TRIANGULAR MATRIX WITH REAL VEC 390

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C          DIAGONAL ELEMENTS.  IN ADDITION, LOCATION BR(N,1) CONTAINS VEC 400
C          THE TOLERANCE QUANTITY (EPSB) COMPUTED AND SAVED IN CQZVAL, VEC 410
C          VEC 420
C          ALFR, ALFI, AND BETA ARE VECTORS WITH COMPONENTS WHOSE VEC 430
C          RATIOS ((ALFR+I*ALFI)/BETA) ARE THE GENERALIZED VEC 440
C          EIGENVALUES.  THEY ARE USUALLY OBTAINED FROM CQZVAL, VEC 450
C          VEC 460
C          Z=(ZR,ZI) CONTAINS THE TRANSFORMATION MATRIX PRODUCED IN THE VEC 470
C          REDUCTIONS BY CQZHEB AND CQZVAL, IF PERFORMED. VEC 480
C          IF THE EIGENVECTORS OF THE TRIANGULAR PROBLEM ARE VEC 490
C          DESIRED, Z MUST CONTAIN THE IDENTITY MATRIX. VEC 500
C          VEC 510
C          ON OUTPUT- VEC 520
C          VEC 530
C          A IS UNALTERED, VEC 540
C          VEC 550
C          B HAS BEEN DESTROYED, VEC 560
C          VEC 570
C          ALFR, ALFI, AND BETA ARE UNALTERED, VEC 580
C          VEC 590
C          Z CONTAINS THE EIGENVECTORS.  EACH EIGENVECTOR IS NORMALIZED VEC 600
C          SO THAT THE MODULUS OF ITS LARGEST COMPONENT IS 1.0 . VEC 610
C          VEC 620
C          QUESTIONS AND COMMENTS SHOULD BE DIRECTED TO B. S. GARBOW, VEC 630
C          APPLIED MATHEMATICS DIVISION, ARGONNE NATIONAL LABORATORY VEC 640
C          VEC 650
C          ----- VEC 660
C          VEC 670
C          IF (N .LE. 1) GO TO 1001 VEC 680
C          EPSB = BR(N,1) VEC 690
C          ***** FOR EN=N STEP -1 UNTIL 2 DO -- ***** VEC 700
C          DO 800 NN = 2, N VEC 710
C             EN = N + 2 - NN VEC 720
C             NA = EN - 1 VEC 730
C             ALMR = ALFR(EN) VEC 740
C             ALMI = ALFI(EN) VEC 750
C             BETM = BETA(EN) VEC 760
C          ***** FOR I=EN-1 STEP -1 UNTIL 1 DO -- ***** VEC 770
C             DO 700 II = 1, NA VEC 780
C                 I = EN - II VEC 790
C                 R = 0.0 VEC 800
C                 RI = 0.0 VEC 810
C                 M = I + 1 VEC 820
C             VEC 830
C             DO 610 J = M, EN VEC 840
C                 T = BETM * AR(I,J) - ALMR * BR(I,J) + ALMI * BI(I,J) VEC 850
C                 TI = BETM * AI(I,J) - ALMR * BI(I,J) - ALMI * BR(I,J) VEC 860
C                 IF (J .EQ. EN) GO TO 605 VEC 870
C                 XI = T * BI(J,EN) + TI * BR(J,EN) VEC 880
C                 T = T * BR(J,EN) - TI * BI(J,EN) VEC 890
C                 TI = XI VEC 900
C             605 R = R + T VEC 910
C                 RI = RI + TI VEC 920
C             610 CONTINUE VEC 930
C             VEC 940
C                 T = ALMR * BETA(I) - BETM * ALFR(I) VEC 950
C                 TI = ALMI * BETA(I) - BETM * ALFI(I) VEC 960
C                 IF (T .EQ. 0.0 .AND. TI .EQ. 0.0) T = EPSB VEC 970
C                 Z3 = CMLPX(R,RI) / CMLPX(T,TI) VEC 980
C                 BR(I,EN) = REAL(Z3) VEC 990
C                 BI(I,EN) = AIMAG(Z3) VEC 1000
C             700 CONTINUE VEC 1010
C             VEC 1020
C             800 CONTINUE VEC 1030
C             ***** END BACK SUBSTITUTION. VEC 1040
C             TRANSFORM TO ORIGINAL COORDINATE SYSTEM. VEC 1050
C             FOR J=N STEP -1 UNTIL 2 DO -- ***** VEC 1060
C             DO 880 JJ = 2, N VEC 1070
C                 J = N + 2 - JJ VEC 1080
C                 M = J - 1 VEC 1090
C             VEC 1100
C             DO 880 I = 1, N VEC 1110
C             VEC 1120
C             DO 860 K = 1, M VEC 1130
C                 ZR(I,J) = ZR(I,J) + ZR(I,K) * BR(K,J) - ZI(I,K) * BI(K,J) VEC 1140
C                 ZI(I,J) = ZI(I,J) + ZR(I,K) * BI(K,J) + ZI(I,K) * BR(K,J) VEC 1150

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860          CONTINUE                                VEC 1160
C
880 CONTINUE                                VEC 1170
C ***** NORMALIZE SO THAT MODULUS OF LARGEST    VEC 1180
C          COMPONENT OF EACH VECTOR IS 1 *****  VEC 1190
C          DO 950 J = 1, N                          VEC 1200
C              T = 0.0                               VEC 1210
C
C          DO 930 I = 1, N                          VEC 1220
C              R = CABS(CMPLX(ZR(I,J),ZI(I,J)))       VEC 1230
C              IF (R .GT. T) T = R                  VEC 1240
C          CONTINUE                                VEC 1250
C
C          DO 940 I = 1, N                          VEC 1260
C              ZR(I,J) = ZR(I,J) / T                VEC 1270
C              ZI(I,J) = ZI(I,J) / T                VEC 1280
C          CONTINUE                                VEC 1290
C
C          DO 950 CONTINUE                          VEC 1300
C
C          DO 940 I = 1, N                          VEC 1310
C              ZR(I,J) = ZR(I,J) / T                VEC 1320
C              ZI(I,J) = ZI(I,J) / T                VEC 1330
C          CONTINUE                                VEC 1340
C
C          DO 950 CONTINUE                          VEC 1350
C
C          DO 1001 RETURN                            VEC 1360
C ***** LAST CARD OF CQZVEC *****             VEC 1370
C          END                                       VEC 1380

C
C          CGG 10
C THIS DRIVER TESTS QZ FOR THE CLASS OF COMPLEX GENERALIZED MATRIXCGG 20
C SYSTEMS EXHIBITING THE USE OF QZ TO FIND ALL THE EIGENVALUES CGG 30
C AND EIGENVECTORS FOR THE EIGENPROBLEM A*X = (LAMBDA)*B*X . CGG 40
C
C          CGG 50
C THE DIMENSION OF A,AI,B,BI,Z, AND ZI SHOULD BE NM BY NM. CGG 60
C THE DIMENSION OF ALFR,ALFI,BETA, AND NORM SHOULD BE NM. CGG 70
C THE DIMENSION OF AHOLD AND BHOLD SHOULD BE NM BY NM. CGG 80
C THE DIMENSION OF AHOLDI AND BHOLDI SHOULD BE NM BY NM. CGG 90
C HERE NM = 20. CGG 100
C
C          CGG 110
C          REAL A( 20, 20),B( 20, 20),Z( 20, 20), CGG 120
C          $ ALFR( 20),ALFI( 20),BETA( 20),NORM( 20), CGG 130
C          $ RESDUL,EPS1, AI(20,20),BI(20,20),ZI(20,20), CGG 140
C          $ AHOLD( 20, 20),BHOLD( 20, 20), AHOLDI(20,20),BHOLDI(20,20)CGG 150
C          REAL AMAX1 CGG 160
C          INTEGER ERROR CGG 170
C          DATA IWRITE/6/ CGG 180
C
C          CGG 190
C          NM = 20 CGG 200
C          10 CALL RMATIN(NM,N,A,B,AHOLD,BHOLD,0) CGG 210
C          CALL RMATIN(NM,N,AI,BI,AHOLDI,BHOLDI,0) CGG 220
C          WRITE(IWRITE,20) CGG 230
C          $ N CGG 240
C          20 FORMAT(30H1THE FULL MATRIX A OF ORDER, CGG 250
C          $ I4,22H IS (PRINTED BY ROWS)/) CGG 260
C          DO 30 I = 1,N CGG 270
C          30 WRITE(IWRITE,40) CGG 280
C          $ (A(I,J),AI(I,J),J=1,N) CGG 290
C          40 FORMAT(5(2F6.0,1H1,3X)) CGG 300
C          WRITE(IWRITE,41) CGG 310
C          $ N CGG 320
C          41 FORMAT(30H0THE FULL MATRIX B OF ORDER, CGG 330
C          $ I4,22H IS (PRINTED BY ROWS)/) CGG 340
C          DO 42 I = 1,N CGG 350
C          42 WRITE(IWRITE,43) CGG 360
C          $ (B(I,J),BI(I,J),J=1,N) CGG 370
C          43 FORMAT(5(2F6.0,1H1,3X)) CGG 380
C
C          CGG 390
C          EIGENVALUES AND EIGENVECTORS USING CQZVAL AND CQZVEC CGG 400
C
C          CGG 410
C          EPS1 = 0.0 CGG 420
C          CALL CQZHES(NM,N,A,AI,B,BI,.TRUE.,Z,ZI) CGG 430
C          CALL CQZVAL(NM,N,A,AI,B,BI,EPS1,ALFR,ALFI,BETA, CGG 440
C          $ .TRUE.,Z,ZI,ERROR) CGG 450
C          CALL CQZVEC(NM,N,A,AI,B,BI,ALFR,ALFI,BETA,Z,ZI) CGG 460
C          WRITE(IWRITE,292) CGG 470
C          292 FORMAT(/15X,7HALFR(I),19X,7HALFI(I),19X,7HBETA(I)) CGG 480
C          DO 295 I = 1,N CGG 490
C          WRITE(IWRITE,293) CGG 500
C          $ I,ALFR(I),ALFI(I),BETA(I) CGG 510

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293     FORMAT(I2,3(1PE23.6,3X))
295     CONTINUE
      CALL  RMATIN(NM,N,A,B,AHOLD,BHOLD,1)
      CALL  RMATIN(NM,N,AI,BI,AHOLDI,BHOLDI,1)
      CALL  RMATIN(NM,N,AHOLD,AHOLDI,Z,ZI,1)
      CALL  CGGWZR(NM,N,A,AI,B,BI,ALFR,ALFI,BETA,AHOLD,AHOLDI,
      $      NORM,RESDUL)
      WRITE(IWRITE,300)
300    FORMAT(/14X,35HCOMPUTED EIGENVALUE AND EIGENVECTOR,20X,8HRESIDUAL)
C
      DO 510 K = 1, N
      BETA(K) = AMAX1(BETA(K),1.0E-50)
      ALFR(K) = ALFR(K) / BETA(K)
      ALFI(K) = ALFI(K) / BETA(K)
C
C     ONE EIGENVECTOR.
C
340    WRITE(IWRITE,350)
      $      K,ALFR(K),ALFI(K),NORM(K),(Z(I,K),
      $      ZI(I,K),I=1,N)
350    FORMAT(/I2,1P2E23.6,E29.2/(5X,2E23.6))
510    CONTINUE
      GO TO 10
      END
      SUBROUTINE RMATIN(NM,N,A,B,AHOLD,BHOLD,INITIL)
C
C     THIS INPUT SUBROUTINE READS TWO REAL MATRICES A AND B FROM
C     SYSIN OF ORDER N.
C     TO GENERATE THE MATRICES INITIALLY, INITIL IS TO BE 0.
C     TO REGENERATE THE MATRICES FOR THE PURPOSE OF THE RESIDUAL
C     CALCULATION, INITIL IS TO BE 1.
C
C     THIS ROUTINE IS CATALOGUED AS EISPDRV4(RSGREADI).
C
      REAL A(NM,NM),B(NM,NM),AHOLD(NM,NM),BHOLD(NM,NM)
      REAL FLOAT
      INTEGER IA( 20), IB( 20)
      DATA IREADA/1/,IREADB/2/,IWRITE/6/
C
      IF( INITIL .EQ. 1 ) GO TO 30
      READ(IREADA,5)
      $      N, M
5     FORMAT(I6,6X,I6)
      IF( N .EQ. 0 ) GO TO 70
      IF( M .NE. 1 ) GO TO 16
      DO 10 I = 1,N
      READ(IREADA,17)
      $      (IA(J), J=I,N)
      DO 9 J = I,N
      A(I,J) = FLOAT(IA(J))
9     A(J,I) = A(I,J)
10    CONTINUE
11    READ(IREADB,5)
      $      N,M
      IF( M .NE. 1 ) GO TO 20
      DO 15 I = 1,N
      READ(IREADB,17)
      $      (IB(J), J=I,N)
      DO 14 J = I,N
      B(I,J)=FLOAT(IB(J))
14    B(J,I)=B(I,J)
15    CONTINUE
      GO TO 22
16    DO 18 I = 1,N
      READ(IREADA,17)
      $      (IA(J), J=1,N)
17    FORMAT(6I12)
      DO 18 J = 1,N
18    A(I,J) = FLOAT(IA(J))
      GO TO 11
20    DO 25 I = 1,N
      READ(IREADB,17)
      $      (IB(J),J=1,N)
      DO 25 J = 1,N
25    B(I,J) = FLOAT(IB(J))
22    DO 23 I = 1,N

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      DO 23 J = 1,N
        BHOLD(I,J) = B(I,J)
23     AHOLD(I,J) = A(I,J)
      RETURN
30 DO 40 I = 1,N
      DO 40 J = 1,N
        B(I,J) = BHOLD(I,J)
40     A(I,J) = AHOLD(I,J)
      RETURN
70 WRITE(IWRITE,80)
80 FORMAT(47HOEND OF DATA FOR SUBROUTINE  RMATIN(RSGREADI). /1H1)
      STOP
      END
C
C -----
C
C SUBROUTINE CGGWZR(NM,N,A,AI,B,BI,ALFR,ALFI,BETA,Z,ZI,NORM,RESDUL)
C
C REAL NORM(N), ALFR(N), ALFI(N), BETA(N), CPART(2), A(NM,N),
C $   B(NM,N), Z(NM,N), XR, XI, S, SUMZ, SUMR, SUMI,
C $   RESDUL,NORMAB,SUMA,SUMB,NORMA,NORMB,
C $   SUMR2,SUMI2,SUMR3,SUMI3, AI(NM,N),BI(NM,N),ZI(NM,N)
C REAL CABS,ABS,FLOAT,AMAX1
C COMPLEX C, C1
C COMPLEX CMLX
C EQUIVALENCE(C1,CPART(1))
C
C THIS SUBROUTINE FORMS THE 1-NORM OF THE RESIDUAL MATRIX
C A*Z-B*Z*DIAG(W) WHERE A AND B ARE COMPLEX GENERAL MATRICES, Z IS
C A MATRIX WHICH CONTAINS THE EIGENVECTORS OF THE EIGENPROBLEM
C A*Z - B*Z*DIAG(W), AND W STANDS FOR A VECTOR OF CORRESPONDING
C EIGENVALUES OF THE EIGENPROBLEM OBTAINED FROM THE VECTORS ALFR,
C ALFI, AND BETA BY THE CORRESPONDENCES
C  $W(J) = (ALFR(J) + I*ALFI(J)) / BETA(J)$  .
C ALL NORMS APPEARING IN THE COMMENTS BELOW ARE 1-NORMS.
C
C INPUT-
C
C NM IS THE ROW DIMENSION OF TWO-DIMENSIONAL ARRAY PARAMETERS
C AS DECLARED IN THE CALLING PROGRAM DIMENSION STATEMENT,
C
C N IS THE ORDER OF THE MATRICES A AND B,
C
C A(NM,N),AI(NM,N),B(NM,N), AND BI(NM,N) ARE ARRAYS WHICH
C CONTAIN THE MATRICES OF THE SYSTEM,
C
C Z(NM,N) AND ZI(NM,N) ARE ARRAYS WHICH CONTAIN THE
C EIGENVECTORS OF THE SYSTEM,
C
C ALFR(N), ALFI(N), AND BETA(N) ARE ARRAYS CONTAINING THE
C COMPONENTS OF THE EIGENVALUES OF THE SYSTEM.
C
C OUTPUT-
C
C Z(NM,N) AND ZI(NM,N) ARE ARRAYS WHICH CONTAIN THE NORMALIZED
C APPROXIMATE EIGENVECTORS OF THE SYSTEM. THE EIGENVECTORS
C ARE NORMALIZED USING THE 1-NORM IN SUCH A WAY THAT THE FIRST
C ELEMENT WHOSE MAGNITUDE IS LARGER THAN THE NORM OF THE
C EIGENVECTOR DIVIDED BY N IS REAL AND POSITIVE,
C
C NORM(N) IS AN ARRAY SUCH THAT FOR EACH K,
C
C 
$$NORM(K) = \frac{..BETA(K)*A*Z(K)-ALFA*B*Z(K)..}{..Z(K)..*(BETA(K)*..A..+.ALFA(K)*..B..)}$$

C
C WHERE Z(K) IS THE K-TH EIGENVECTOR AND ALFA = (ALFR + I*ALFI),
C
C RESDUL IS THE REAL NUMBER
C 
$$..BETA*A*Z-ALFA*B*Z../(..Z..*(BETA*..A..+.ALFA*..B..))$$

C -----
C
C NORMB = 0.0
C NORMA = 0.0
C RESDUL = 0.0

```

CGG 1280
CGG 1290
CGG 1300
CGG 1310
CGG 1320
CGG 1330
CGG 1340
CGG 1350
CGG 1360
CGG 1370
CGG 1380
CGG 1390
CGG 1400
CGG 1410
CGG 1420
CGG 1430
CGG 1440
CGG 1450
CGG 1460
CGG 1470
CGG 1480
CGG 1490
CGG 1500
CGG 1510
CGG 1520
CGG 1530
CGG 1540
CGG 1550
CGG 1560
CGG 1570
CGG 1580
CGG 1590
CGG 1600
CGG 1610
CGG 1620
CGG 1630
CGG 1640
CGG 1650
CGG 1660
CGG 1670
CGG 1680
CGG 1690
CGG 1700
CGG 1710
CGG 1720
CGG 1730
CGG 1740
CGG 1750
CGG 1760
CGG 1770
CGG 1780
CGG 1790
CGG 1800
CGG 1810
CGG 1820
CGG 1830
CGG 1840
CGG 1850
CGG 1860
CGG 1870
CGG 1880
CGG 1890
CGG 1900
CGG 1910
CGG 1920
CGG 1930
CGG 1940
CGG 1950
CGG 1960
CGG 1970
CGG 1980
CGG 1990
CGG 2000
CGG 2010
CGG 2020
CGG 2030

```

C
DO 20 I = 1,N
  SUMA = 0.0
  SUMB = 0.0
  DO 10 L = 1,N
    SUMA = SUMA + ABS(A(L,I)) + ABS(AI(L,I))
10    SUMB = SUMB + ABS(B(L,I)) + ABS(BI(L,I))
    NORMA = AMAX1(NORMA,SUMA)
20    NORMB = AMAX1(NORMB,SUMB)
C
DO 160 I=1,N
  S = 0.0
  SUMZ = 0.0
  DO 110 L = 1,N
    SUMR = 0.0
    SUMI = 0.0
    SUMR2 = 0.0
    SUMI2 = 0.0
    SUMZ = SUMZ + CABS(CMPLX(Z(L,I),ZI(L,I)))
C
DO 100 K=1,N
  SUMR = SUMR + B(L,K)*Z(K,I) - BI(L,K)*ZI(K,I)
  SUMI = SUMI + B(L,K)*ZI(K,I) + BI(L,K)*Z(K,I)
  SUMR2 = SUMR2 + A(L,K)*Z(K,I) - AI(L,K)*ZI(K,I)
100  SUMI2 = SUMI2 + A(L,K)*ZI(K,I) + AI(L,K)*Z(K,I)
  SUMR3 = -ALFR(I)*SUMR + ALFI(I)*SUMI
  SUMI3 = -ALFI(I)*SUMR - ALFR(I)*SUMI
C
110  S = S+CABS(CMPLX(SUMR3,SUMI3)+CMPLX(SUMR2,SUMI2)*BETA(I))
  NORMAB = NORMA*BETA(I)+NORMB*CABS(CMPLX(ALFR(I),ALFI(I)))
  IF( NORMAB .EQ. 0.0 ) NORMAB = 1.0
C
  NORM(I) = SUMZ
  IF( SUMZ .EQ. 0.0 ) GO TO 150
C
  *****THIS LOOP WILL NEVER BE COMPLETED SINCE THERE WILL
C
  ALWAYS EXIST AN ELEMENT IN THE VECTOR (Z(I),ZI(I))
C
  LARGER THAN ..(Z(I),ZI(I)).. /N*****
C
DO 120 L=1,N
  IF(CABS(CMPLX(Z(L,I),ZI(L,I))) .GE. NORM(I)/FLOAT(N))
  $ GO TO 130
120  CONTINUE
C
130  XR = NORM(I)*Z(L,I)/CABS(CMPLX(Z(L,I),ZI(L,I)))
  XI = NORM(I)*ZI(L,I)/CABS(CMPLX(Z(L,I),ZI(L,I)))
  C = CMPLX(XR,XI)
C
DO 140 L= 1,N
  C1 = CMPLX(Z(L,I),ZI(L,I))/C
  Z(L,I) = CPART(1)
140  ZI(L,I) = CPART(2)
C
  NORM(I) = S/(NORM(I)*NORMAB)
150  RESDUL = AMAX1(NORM(I),RESDUL)
160  CONTINUE
C
  RETURN
  END

```

CGG 2040
CGG 2050
CGG 2060
CGG 2070
CGG 2080
CGG 2090
CGG 2100
CGG 2110
CGG 2120
CGG 2130
CGG 2140
CGG 2150
CGG 2160
CGG 2170
CGG 2180
CGG 2190
CGG 2200
CGG 2210
CGG 2220
CGG 2230
CGG 2240
CGG 2250
CGG 2260
CGG 2270
CGG 2280
CGG 2290
CGG 2300
CGG 2310
CGG 2320
CGG 2330
CGG 2340
CGG 2350
CGG 2360
CGG 2370
CGG 2380
CGG 2390
CGG 2400
CGG 2410
CGG 2420
CGG 2430
CGG 2440
CGG 2450
CGG 2460
CGG 2470
CGG 2480
CGG 2490
CGG 2500
CGG 2510
CGG 2520
CGG 2530
CGG 2540
CGG 2550
CGG 2560
CGG 2570
CGG 2580
CGG 2590
CGG 2600

```

5
-238      86      164      -166      56      DATA1 1
  76      -96      40      60      -60      DATA1 2
  118      55      -13      34      -176     DATA1 3
 -314      132     114      -90      -424     DATA1 4
  -54      -205     109      158      -38      DATA1 5
5
-344      178      240      -308     158      DATA1 7
  152      -128     -32      184      -136     DATA1 8
  284      -182     460      -192     -214     DATA1 9
 -160      78      296      -164     -374     DATA1 10
  -24      -400     148      312      -96      DATA1 11
5
  41      -143     -20      20      104      DATA1 12
  148      144      -6      -78      8        DATA1 13
  -19      87      4        -56     -164     DATA1 14
 -60      -81      99      34      84       DATA1 15
  1        133     132     -46     -12      DATA1 16
5

```

DATA1 1
DATA1 2
DATA1 3
DATA1 4
DATA1 5
DATA1 6
DATA1 7
DATA1 8
DATA1 9
DATA1 10
DATA1 11
DATA1 12
DATA1 13
DATA1 14
DATA1 15
DATA1 16
DATA1 17
DATA1 18
DATA1 19

	-369	-747	-1368	486	-432	DATA1 20
	261	666	-1152	45	-540	DATA1 21
	819	243	1548	-954	180	DATA1 22
	-945	-279	171	441	-144	DATA1 23
	-468	747	774	-45	-216	DATA1 24
5						DATA1 25
	-15	-143	-83	41	55	DATA1 26
	100	144	-60	-60	-34	DATA1 27
	-19	87	4	-56	-164	DATA1 28
	-116	-81	36	55	35	DATA1 29
	-39	133	87	-31	-47	DATA1 30
5						DATA1 31
	-1635	-173	-1142	1012	-566	DATA1 32
	-829	1128	-1050	495	-710	DATA1 33
	971	187	1558	-1016	218	DATA1 34
	-562	-944	-1310	238	-937	DATA1 35
	356	-149	-875	-434	-1015	DATA1 36
0						DATA1 37
5						DATA2 1
	388	-386	-250	556	-396	DATA2 2
	-304	384	-160	-240	240	DATA2 3
	-658	-73	-109	-118	406	DATA2 4
	-640	204	-692	288	-192	DATA2 5
	-162	631	131	-758	278	DATA2 6
5						DATA2 7
	94	-122	-14	130	-62	DATA2 8
	-76	64	16	-92	68	DATA2 9
	-136	100	-250	100	96	DATA2 10
	-10	-42	-90	66	154	DATA2 11
	-72	158	52	-184	76	DATA2 12
5						DATA2 13
	90	180	36	-90	-72	DATA2 14
	-105	-210	-42	105	84	DATA2 15
	-90	-180	-36	90	72	DATA2 16
	75	150	30	-75	-60	DATA2 17
	-75	-150	-30	75	60	DATA2 18
5						DATA2 19
	161	335	182	-162	-36	DATA2 20
	-169	-322	24	167	204	DATA2 21
	-211	-307	-160	186	36	DATA2 22
	205	215	45	-165	-80	DATA2 23
	-48	-299	-102	89	88	DATA2 24
5						DATA2 25
	90	180	36	-90	-72	DATA2 26
	-105	-210	-42	105	84	DATA2 27
	-90	-180	-36	90	72	DATA2 28
	75	150	30	-75	-60	DATA2 29
	-75	-150	-30	75	60	DATA2 30
5						DATA2 31
	307	253	163	-235	-39	DATA2 32
	-49	-410	46	85	182	DATA2 33
	-229	-253	-200	234	84	DATA2 34
	171	301	184	-130	25	DATA2 35
	-134	-185	59	146	195	DATA2 36

ACM Transactions on Mathematical Software, Vol. 8, No. 4, December 1982, Page 402.

Remark on Algorithm 535

The QZ Algorithm to Solve the Generalized Eigenvalue Problem [B.S. Garbow, *ACM Trans. Math. Softw.* 4, 4 (Dec. 1978), 404-410]

B.S. Garbow [Received 10 June 1982; accepted 10 June 1982]

Applied Mathematics Division, Argonne National Laboratory, 9700 S. Cass Avenue, Argonne, IL 60439.

Replace the single card with sequence label **VAL 153**:

IF (BI (L, L) .EQ. 0.0) GO TO 98

with

IF (B11 .EQ. 0.0) GO TO 98

The error that this change corrects results occasionally in eigenvalues with the wrong sign.

ALGORITHM 536

An Efficient One-Way Enciphering Algorithm [Z]

H. D. KNOBLE
The Pennsylvania State University

Key Words and Phrases: one-way security transformation, password, encipher, decipher, multiprecision integer arithmetic
CR Categories: 3.15, 4.49
Language: Fortran

DESCRIPTION

This algorithm is a Fortran implementation of the procedures developed in [1]. Its purpose is to serve as a machine independent model for studying the evaluation of polynomials mod P and for the implementation of more efficient machine dependent system utility programs for enciphering passwords.

REFERENCES

1. KNOBLE, H.D., FORNEY, C., AND BADER, F.S. An efficient one-way enciphering algorithm. *ACM Trans. Math. Software* 5, 1 (March 1979), 97-107.

ALGORITHM

```

C----- SAMPLE PROGRAM TO ENCIIPHER 72-BIT PASSWORDS ACROSS      00000010
C MACHINES WITH 3 DIFFERENT WORD SIZES. OUTPUT IS TABLE 2A OR 2B. 00000020
C TO RUN THIS PROGRAM COMMENTS BEGINNING C32, C36, OR C60 MUST BE 00000030
C ACTIVATED BY REPLACING THESE FIRST THREE CHARACTERS WITH BLANKS. 00000040
C TABLE 2A WILL BE PRODUCED ON IBM 360/370 OR XEROX SIGMA SERIES    00000050
C 32-BIT MACHINES BY ACTIVATING COMMENTS BEGINNING WITH C32.      00000060
C TABLE 2B WILL BE PRODUCED ON UNIVAC 1100, HONEYWELL 600/6000, AND 00000070
C DEC SYSTEM 10/20 36-BIT MACHINES WHEN ACTIVATING COMMENTS BEGINNING 00000080
C WITH C36. TABLE 2B WILL BE PRODUCED ON THE CDC 6600/7600 60-BIT 00000090
C MACHINES WHEN ACTIVATING COMMENTS BEGINNING WITH C60.           00000100
C                                                                    00000110
C INPUT TO THIS PROGRAM ARE THE FOLLOWING 10 CARDS BEGINNING WITH C= 00000120
C= HEX(000000000000000000) OCTAL(000000000000000000000000) DECIMAL(0) 00000130
C= HEX(000000000000000001) OCTAL(0000000000000000000001) DECIMAL(1) 00000140
C= HEX(000000000000000002) OCTAL(0000000000000000000002) DECIMAL(2) 00000150
C= HEX(000000000000000003) OCTAL(0000000000000000000003) DECIMAL(3) 00000160
C= HEX(FFFFFFFFFFFFFFFFFA3) OCTAL(7777777777777777777777643) DECIMAL(P) 00000170
C= HEX(FFFFFFFFFFFFFFFFFA4) OCTAL(7777777777777777777777644) DECIMAL(P+1) 00000180
C= HEX(FFFFFFFFFFFFFFFFFA5) OCTAL(7777777777777777777777645) DECIMAL(P+2) 00000190
C= HEX(00000000000000005C) OCTAL(0000000000000000000000134) DECIMAL(92) 00000200
C= HEX(FFFFFFFFFFFFFFFFFFF) OCTAL(7777777777777777777777777) DECIMAL(P+92) 00000210
C= HEX(555555555555555555) OCTAL(25252525252525252525252525) BINARY(01...)00000220

```

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ACM Transactions on Mathematical Software, Vol. 5, No. 1, March 1979, Pages 108-111.

```

C
C-----
C      INTEGER A,NN,N(6)
C32  INTEGER X(3),FX(3),P(3),AC(3,6),WORK(43),NP,IW
C36  INTEGER X(2),FX(2),P(2),AC(2,6),WORK(31),NP,IW
C60  INTEGER X(2),FX(2),P(2),AC(2,6),WORK(31),NP,IW
C-----A IS CHOSEN TO MAKE P=2**72-A A PRIME INTEGER IN THE
C      INTERVAL (2**(NP*M)-2*(M-1), 2**(NP*M)). NP*M=72 HERE.
      DATA A/93/, NN/6/
      DATA N(1)/1048569/,N(2)/524285/,N(3)/3/,N(4)/2/,N(5)/1/,N(6)/0/
C.32.....IBM, XEROX 32-BIT MACHINES, HEXADECIMAL CONSTANTS(Z).
C32  DATA NP/3/, IW/43/
C32  DATA AC(1,1),AC(2,1),AC(3,1)/Z00000000,Z00000000,Z00000001/,
C32  * AC(1,2),AC(2,2),AC(3,2)/Z00FFFFFF,Z00FFFFFF,Z00FFFFFFA1/,
C32  * AC(1,3),AC(2,3),AC(3,3)/Z00FFFFFF,Z00FFFFFF,Z00FFFFFF83/,
C32  * AC(1,4),AC(2,4),AC(3,4)/Z00FFFFFF,Z00FFFFFF,Z00FFFFFF7B/,
C32  * AC(1,5),AC(2,5),AC(3,5)/Z00FFFFFF,Z00FFFFFF,Z00FFFFFF9E/,
C32  * AC(1,6),AC(2,6),AC(3,6)/Z00FFFFFF,Z00FFFFFF,Z00FFFFFF02/
C
C.36.....UNIVAC, HONEYWELL, DEC 36-BIT MACHINES, OCTAL CONSTANTS(O).
C36  DATA NP/2/, IW/31/
C36  DATA AC(1,1),AC(2,1)/00000000000000000000000000000001/,
C36  * AC(1,2),AC(2,2)/0777777777777777,07777777777641/,
C36  * AC(1,3),AC(2,3)/0777777777777777,07777777777603/,
C36  * AC(1,4),AC(2,4)/0777777777777777,07777777777573/,
C36  * AC(1,5),AC(2,5)/0777777777777777,07777777777636/,
C36  * AC(1,6),AC(2,6)/0777777777777777,07777777777402/
C
C.60.....CDC 6600/7600 60-BIT +, 48-BIT */ , OCTAL CONSTANTS(B).
C60  DATA NP/2/, IW/31/
C60  DATA AC(1,1),AC(2,1)/00000000000000000000000000000001B/,
C60  * AC(1,2),AC(2,2)/0000000007777777777B,00000000077777777641B/,
C60  * AC(1,3),AC(2,3)/0000000007777777777B,00000000077777777603B/,
C60  * AC(1,4),AC(2,4)/0000000007777777777B,00000000077777777573B/,
C60  * AC(1,5),AC(2,5)/0000000007777777777B,00000000077777777636B/,
C60  * AC(1,6),AC(2,6)/0000000007777777777B,00000000077777777402B/
C
      WRITE(6,1)
C32 1 FORMAT(45H1HEXADECIMAL INPUT(X) HEXADECIMAL OUTPUT(FX)/1X)
C36 1 FORMAT(15H1OCTAL INPUT(X),14X,16HOCTAL OUTPUT(FX)/1X)
C60 1 FORMAT(15H1OCTAL INPUT(X),14X,16HOCTAL OUTPUT(FX)/1X)
C-----INPUT AND ENCIIPHER 10 BIT STRINGS.
      DO 4 I=1,10
      READ(5,2) X
C32 2  FORMAT(7X,3Z6)
C36 2  FORMAT(33X,2012)
C60 2  FORMAT(33X,2012)
      CALL PURDY(X,FX,A,P,N,AC,WORK,NP,NN,IW,NP,IERR)
      IF(IERR.NE.0) GO TO 5
      WRITE(6,3) X,FX
C32 3  FORMAT(1X,3Z6,4X,3Z6)
C36 3  FORMAT(1X,2012,4X,2012)
C60 3  FORMAT(1X,2012,4X,2012)
      4 CONTINUE
      STOP
      5 WRITE(6,6) IERR
      6 FORMAT(46H0***ERROR RETURN FROM SUBROUTINE PURDY. IERR=,I2)
      STOP
      END
      BLOCK DATA
      COMMON/MACHIN/ ISIZE,M,MAXPOS,MINNEG,MASK
C32  DATA ISIZE/32/,M/24/,MAXPOS/Z007FFFFFF/,MINNEG/Z00800000/,
C32  * MASK/Z01000000/
C36  DATA ISIZE/36/,M/36/,MAXPOS/0377777777777777/,MINNEG/0400000000000000/,
C36  * MASK/0/
C60  DATA ISIZE/60/,M/36/,MAXPOS/00000000377777777777B/,
C60  * MINNEG/0000000040000000000000B/,MASK/00000001000000000000B/
      END
      SUBROUTINE PURDY(X,FX,A,PRIME,N,ACOEFF,WORK,NP,NN,NW,IDIMA,IERR)
C =====ROUTINE TO EVALUATE PURDY'S IRREVERSIBLE ENCIIPHERING FUNCTION...
C      FX=F(X)=P(X) MOD PRIME WHERE
C      P(X)= SUM(ACOEFF(I)*X**N(I)), I=1,2,3,...,NN.

```



```

C   ACOEFF IS GIVEN AS AN IDIMA BY NN ARRAY REPRESENTING THE UNSIGNED 00000820
C   MULTIPRECISION INTEGER COEFFICIENTS OF THE ENCIPHERING 00000830
C   POLYNOMIAL F(X) ABOVE. THESE COEFFICIENTS SHOULD BE 00000840
C   CHOSEN TO HAVE (LARGE) VALUES, LESS THAN 2**Q-A. THE 00000850
C   PRECISION OF THE ELEMENTS OF ACOEFF IS NP AND THEREFORE 00000860
C   IDIMA USUALLY WOULD BE CHOSEN SUCH THAT IDIMA=NP. 00000870
C   00000880
C   WORK IS GIVEN AS AN WORKSPACE ARRAY OF NW=12*NP+7 INTEGER WORDS. 00000890
C   00000900
C   NP IS GIVEN AS AN INTEGER WHOSE VALUE IS THE PRECISION OF 00000910
C   MULTIPRECISION INTEGERS. THAT IS, NP IS THE NUMBER OF 00000920
C   INTEGER COMPUTER WORDS NECESSARY TO STORE THE UNSIGNED 00000930
C   MULTIPRECISION INTEGER 2**Q-1. NP MUST BE GREATER THAN 1. 00000940
C   00000950
C   NN IS GIVEN AS AN INTEGER WHOSE VALUE IS THE NUMBER OF TERMS IN 00000960
C   P(X) ABOVE. NN MUST BE GREATER THAN 0. 00000970
C   00000980
C   NW IS GIVEN AS AN INTEGER WHOSE VALUE IS GREATER THAN OR EQUAL 00000990
C   TO 12*NP+7. NW IS THE DIMENSION OF THE WORKSPACE ARRAY, WORK. 00001000
C   00001010
C   IDIMA IS GIVEN AS AN INTEGER WHOSE VALUE IS EQUAL TO THE 00001020
C   FIRST DIMENSION CONSTANT OF THE ARGUMENT IN THE CALLING 00001030
C   PROGRAM CORRESPONDING TO THE PARAMETER, ACOEFF. NORMALLY 00001040
C   IDIMA=NP. 00001050
C   00001060
C   IERR IS RETURNED AS AN ERROR INDICATOR AS FOLLOWS... 00001070
C   IERR=1 IF NP .LE. 1 OR NN IS NOT POSITIVE, IF NW IS LESS THAN 12*NP+7 00001080
C   OR IF ANY OF THE NN ELEMENTS OF N ARE NEGATIVE. 00001090
C   IERR=2 IF A IS NOT IN ITS RANGE DESCRIBED ABOVE. 00001100
C   IERR=3 IF M IS UNEVEN OR GREATER THAN ISIZE OR IF THE HOST MACHINE 00001110
C   PERFORMS SIGNED MAGNITUDE ARITHMETIC AND M IS UNEVEN OR 00001120
C   GREATER THAN ISIZE-2. 00001130
C   IERR=4 IF THE LEADING ISIZE-M BITS OF ANY ELEMENT OF X ARE NOT ZERO. 00001140
C   IERR=0 OTHERWISE. FX AND PRIME ARE RETURNED UNCHANGED IF IERR 00001150
C   IS RETURNED NON-ZERO. 00001160
C   00001170
C   NOTES: THIS ALGORITHM IS PRESENTED HERE AS A GENERAL EVALUATION 00001180
C   SCHEME AND THE FORTRAN CODE IS INTENDED TO SERVE AS A MODEL 00001190
C   FOR CONSTRUCTING A SIMILAR SYSTEM UTILITY FUNCTION. 00001200
C   TO PROGRAM AN EFFICIENT RE-ENTERABLE ASSEMBLER VERSION OF 00001210
C   THIS ALGORITHM TO BE USED AS PART OF A SYSTEM PASSWORD 00001220
C   AUTHORIZATION SCHEME THE FOLLOWING IS RECOMMENDED... 00001230
C   00001240
C   ---- CHOOSE M=ISIZE (EVEN FOR SIGNED MAGNITUDE MACHINES) AND 00001250
C   USE EQUIVALENT MACHINE INSTRUCTIONS IN PLACE OF SUBROUTINES 00001260
C   LADD, LSUB, AND KOMP. ELIMINATE SUBROUTINES SPLIT, JOIN, 00001270
C   ISIGNM AND ALL REFERENCES TO COMMON BLOCK /MACHIN/. 00001280
C   IN SUBROUTINE MPMLT CUT THE PRECISION FROM 2*NP TO NP BY USING 00001290
C   THE EQUIVALENT OF A LOGICAL MULTIPLY MACHINE INSTRUCTION. 00001300
C   00001310
C   ---- CHOOSE A FAMILY OF ENCIPHERING POLYNOMIALS AND FACTOR IT 00001320
C   BEFORE EVALUATION. FOR EXAMPLE, IF N(1)=2*N(2)-1, N(3)=3, 00001330
C   N(4)=2, N(5)=1, N(6)=0, NN=6, THEN P(X) FACTORS TO... 00001340
C   X**(N(2)-1)*(X*(X**(N(2)-1)+ACOEFF(1,1)) + 00001350
C   X*(ACOEFF(1,4)+X*(ACOEFF(1,3)+ACOEFF(1,2)*X)). 00001360
C   THIS REDUCES THE NUMBER OF CALLS TO EXPP THUS SHARPLY REDUCING 00001370
C   COMPUTING TIME. 00001380
C   00001390
C   ---- WHEN PRECISION IS CONSTANT, WORKSPACE AND RELATED PARAMETERS 00001400
C   BECOME CONSTANTS AND THE PROGRAM IS CONSIDERABLY SIMPLIFIED. 00001410
C   ONCE THIS IS DONE PASS ALL WORKSPACE VARIABLES IN COMMON AND 00001420
C   REMOVE ALL DIMENSION PARAMETERS AND WORKSPACE PARAMETERS FROM 00001430
C   THE ARGUMENT/PARAMETER LISTS. THIS SUBSTANTIALLY REDUCES 00001440
C   LINKAGE TIME. 00001450
C   00001460
C   WRITTEN BY H. D. KNOBLE, PENN STATE UNIVERSITY COMPUTATION CENTER, 00001470
C   JUNE 1977. 00001480
C=====00001490
C   INTEGER X(NP),FX(NP),Q,A,PRIME(NP),N(NN),ACOEFF(IDIMA,NN),WORK(NW)00001500
C   COMMON/MACHIN/ISIZE,M,MAXPOS,MINNEG,MASK 00001510
C-----CHECK PARAMETERS FOR DOMAIN ERRORS. 00001520
C   IERR=0 00001530
C   IF(NP.LE.1.OR.NN.LE.0) IERR=1 00001540
C   IF(IDIMA.LT.NP.OR.NW.LT.12*NP+7) IERR=1 00001550
C   IF(A.LE.0.OR.A.GT.MAXPOS) IERR=2 00001560
C   DO 1 I=1,NN 00001570
C   IF(N(I).LT.0) IERR=1 00001580

```

```

1      CONTINUE                                00001590
C-----CHECK HOST MACHINE ASSUMPTIONS ASSUMING MINNEG IS CORRECT. 00001600
      IF((MOD(M,2).NE.0).OR.(M.GT.ISIZE).OR. 00001610
      * (M.GT.ISIZE-2.AND.MINNEG.EQ.0)) IERR=3 00001620
C-----CHECK IF THE LEADING (ISIZE-M) BITS OF THE WORDS OF X 00001630
C      ARE NON-ZERO.                                00001640
      IF(ISIZE.EQ.M) GO TO 3                      00001650
      DO 2 I=1,NP                                00001660
      IF(X(I).GE.MASK) IERR=4                    00001670
2      CONTINUE                                00001680
3      IF(IERR.NE.0) RETURN                      00001690
C-----FOR Q=NP*M, COMPUTE P=2**Q-A.              00001700
      NP1=NP+1                                   00001710
      NP2=NP*2                                   00001720
      DO 4 I=1,NP2                                00001730
4      WORK(I)=0                                  00001740
      WORK(NP2)=A                                 00001750
      CALL MPSTB(WORK,WORK(NP1),PRIME,NP)        00001760
C-----INITIALIZE                                00001770
      N2=NP2+1                                   00001780
      IWE=10*NP+7                               00001790
      IWM=8*NP+7                                 00001800
      IWA=3*NP+2                                00001810
C-----EVALUATE F(X) ONE TERM AT A TIME.          00001820
      DO 5 I=1,NN                                00001830
      CALL EXPP(X,N(I),WORK(NP1),PRIME,WORK(N2),NP,IWE) 00001840
      CALL MULTP(ACOEFF(1,I),WORK(NP1),WORK(NP1),PRIME,WORK(N2),NP,IWM) 00001850
5      CALL ADDP(WORK,WORK(NP1),WORK,PRIME,WORK(N2),NP,IWA) 00001860
C-----SET RESULT = F(X). FX AND X MAY BE SAME LOCATIONS. 00001870
      DO 6 I=1,NP                                00001880
6      FX(I)=WORK(I)                             00001890
      RETURN                                     00001900
      END                                         00001910

      SUBROUTINE EXPP(X,K,Y,P,WORK,NP,NW)          00001920
C=====COMPUTE Y=X**K MOD P FOR P=2**Q-A FOR UNSIGNED 00001930
C      MULTIPRECISION INTEGERS,X,Y,P OF PRECISION NP. WORK IS A 00001940
C      WORKSPACE ARRAY OF NW=NP*10+7 WORDS.      00001950
C      SEE KNUTH ALGORITHM 4.6.3-A.              00001960
C      INTEGER X(NP),Y(NP),P(NP),WORK(NW)        00001970
      NPP1=NP+1                                   00001980
      IWM=NP*8+7                                 00001990
      N=K                                         00002000
C-----Y=1, Z=X.                                00002010
      DO 1 I=1,NP                                00002020
      Y(I)=0                                      00002030
1      WORK(I)=X(I)                              00002040
      Y(NP)=1                                    00002050
C-----X**0=1 FOR X.GE.0.                        00002060
      IF(N.LT.1) RETURN                          00002070
2      L=N                                        00002080
      N=N/2                                       00002090
      IF(MOD(L,2).EQ.0) GO TO 5                  00002100
      CALL MULTP(Y,WORK,Y,P,WORK(NPP1),NP,IWM)  00002110
      IF(N.EQ.0) RETURN                          00002120
5      CALL MULTP(WORK,WORK,WORK,P,WORK(NPP1),NP,IWM) 00002130
      GO TO 2                                    00002140
      END                                         00002150
      END                                         00002160

      SUBROUTINE MULTP(R,S,RS,P,WORK,NP,NW)       00002170
C=====COMPUTE RS=R*S MOD P FOR P=2**Q-A FOR UNSIGNED 00002180
C      MULTIPRECISION INTEGERS R,S,RS,P OF PRECISION NP. WORK 00002190
C      IS A WORKSPACE ARRAY OF NW=NP*8+7 WORDS.  00002200
C      INTEGER R(NP),S(NP),RS(NP),P(NP),WORK(NW) 00002210
      NP2=2*NP                                   00002220
      IW1=NP2+1                                  00002230
      IW2=IW1+NP2                                00002240
      IW3=IW2+NP2                                00002250
      CALL MPMLT(R,S,WORK,WORK(IW1),WORK(IW2),WORK(IW3),NP,NP, 00002260
      * NP2,NP2,NP2,2*NP2)                      00002270
      CALL MOD2Q(WORK,RS,P,WORK(IW1),NP2,NP,NP*6+7) 00002280
      RETURN                                     00002290
      END                                         00002300

```



```

      RETURN
C-----Y=X-P
2    CALL MPSUB(X,P,Y,NP)
      RETURN
      END
00003010
00003020
00003030
00003040
00003050

      SUBROUTINE MPMLT(IX,IY,IZ,U,V,W,NX,MY,NPM,N2,M2,NPM2)
C=====UNSIGNED MULTIPRECISION INTEGER MULTIPLICATION
C      IZ=IX*IY.  NX,MX,NPM ARE THE PRECISIONS OF THE UNSIGNED
C      INTEGERS IX,IY,IZ RESPECTIVELY.  U,V,W ARE
C      WORKSPACES OF N2=NX*2, M2=MY*2, AND MPN2=(NX+MY)*2 WORDS
C      RESPECTIVELY.  NO INTEGER OVERFLOWS ARE GENERATED.
C
C      SEE KNUTH'S ALGORITHM 4.3.1 M.
C
C      INTEGER IX(NX),IY(MY),IZ(NPM),U(N2),V(M2),W(NPM2)
C      INTEGER T(2),K,SHIFT1,V0
C      COMMON /MACHIN/ ISIZE,M,MAXPOS,MINNEG,MASK
C      SHIFT1=2**(M/2-1)
C-----DOUBLE THE PRECISION SINCE FORTRAN CANNOT ACCESS THE
C      PRODUCT OF TWO M-BIT WORDS.
C      DO 10 I=1,NX
C      I2=I*2
10   CALL SPLIT(IX(I),U(I2-1),U(I2))
C      DO 11 I=1,MY
C      I2=I*2
11   CALL SPLIT(IY(I),V(I2-1),V(I2))
C-----BEGIN KNUTH'S ALGORITHM.
C      DO 1 I=1,NPM2
1    W(I)=0
C      J=M2
C-----THE PROBABILITY OF V(J)=0 IS SMALL.  SKIP STEP 2.
3    I=N2
C      K=0
4    IPJ=I+J
C-----COMPUTE T=U(I)*V(J)+W(I+J)+K WITHOUT INTEGER OVERFLOW.
C      V0=MOD(V(J),SHIFT1)
C      CALL LADD(U(I)*V0,(V(J)-V0)*U(I),IZ)
C      CALL LADD(IZ(2),W(IPJ)+K,T)
C-----COMPUTE K = FLOOR(T/2**(M/2)), W(IPJ)=T MOD 2**(M/2),
C      THE HIGH AND LOW-ORDER HALFS OF T(2).
C      CALL SPLIT(T(2),K,W(IPJ))
5    I=I-1
C      IF(I.GT.0) GO TO 4
C      W(J)=K
6    J=J-1
C      IF(J.GT.0) GO TO 3
C-----CONVERT THE RADIX 2**(M/2) DIGITS TO RADIX 2**M DIGITS.
C      DO 12 I=1,NPM
C      I2=I*2
12   IZ(I)=JOIN(W(I2-1),W(I2))
      RETURN
      END
00003190
00003199
00003200
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00003250
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00003520

      SUBROUTINE MPADD(U,V,W,N,NP1)
C=====UNSIGNED MULTIPRECISION INTEGER ADDITION,
C      W=U+V.  N IS THE PRECISION OF UNSIGNED NUMBERS, IX, IY.
C      W IS THE UNSIGNED NP1-WORD PRODUCT.
C
C      SEE KNUTH'S ALGORITHM 4.3.1-A.
C      INTEGER U(N),V(N),W(NP1),SUM1(2),SUM2(2)
1    J=N
C      K=0
2    CALL LADD(U(J),V(J),SUM1)
C      CALL LADD(SUM1(2),K,SUM2)
C      K=SUM1(1)+SUM2(1)
C      W(J+1)=SUM2(2)
3    J=J-1
C      IF(J.GT.0) GO TO 2
C      W(1)=K
      RETURN
      END
00003530
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00003560
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00003690
00003700

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SUBROUTINE MPSUB(U,V,W,N)                                00003710
C=====UNSIGNED MULTIPRECISION INTEGER SUBTRACTION,    00003720
C      W=U-V IGNORING POSSIBLE BORROW IF U IS LESS THAN V. 00003730
C      N IS THE PRECISION OF UNSIGNED NUMBERS U,V,W.      00003740
C                                                         00003750
C      SEE KNUTH'S ALGORITHM 4.3.1-S.                     00003760
C      INTEGER U(N),V(N),W(N),DIF1(2),DIF2(2)             00003770
1      J=N                                                 00003780
      K=0                                                  00003790
2      CALL LSUB(U(J),V(J),DIF1)                          00003800
      CALL LSUB(DIF1(2),K,DIF2)                           00003810
      K=DIF1(1)+DIF2(1)                                    00003820
      W(J)=DIF2(2)                                         00003830
3      J=J-1                                               00003840
      IF(J.GT.0) GO TO 2                                    00003850
      RETURN                                               00003860
      END                                                  00003870

FUNCTION KOMP(U,V,N)                                     00003880
CXXXXXXXXXX FUNCTION TO COMPARE UNSIGNED MULTIPRECISION 00003890
C      INTEGERS U,V OF PRECISION N. KOMP=-1 IF U IS LESS THAN V, 00003900
C      +1 IF U IS GREATER THAN V, AND 0 IF U=V.           00003910
C                                                         00003920
C      SEE KNUTH PROBLEM 4.3.1-11                          00003930
C      INTEGER U(N),V(N)                                    00003940
      DO 1 J=1,N                                           00003950
C-----IF HIGH-ORDER BITS DIFFER U .NE. V.              00003960
      IF(ISIGN(1,U(J)).NE.ISIGN(1,V(J))) GO TO 2           00003970
C-----ELSE IF U IS LESS THAN V, THEN KOMP=-1           00003980
      IF(U(J).LT.V(J)) GO TO 3                             00003990
C-----ELSE IF U IS GREATER THAN V THEN KOMP=+1.        00004000
      IF(U(J).GT.V(J)) GO TO 4                             00004010
1      CONTINUE                                           00004020
C-----U=V.                                              00004030
      KOMP=0                                               00004040
      RETURN                                               00004050
C-----THE OPERAND WITH HIGH-ORDER BIT ON IS GREATER.   00004060
2      IF(U(J).LT.V(J))GO TO 4                             00004070
3      KOMP=-1                                             00004080
      RETURN                                               00004090
4      KOMP=+1                                             00004100
      RETURN                                               00004110
      END                                                  00004120

FUNCTION JOIN(HALF1,HALF2)                              00004130
CXXXXXXXXXXXXX ROUTINE TO CONVERT TWO RADIX 2**(M/2) DIGITS HALF1, 00004140
C      HALF2 TO A RADIX 2**M DIGIT, JOIN. THAT IS THIS   00004150
C      ROUTINE FORMS A M-BIT NUMBER WITH THE HIGH-ORDER ISIZE-M 00004160
C      BITS EQUAL TO ZERO, THE M/2 HIGH-ORDER BITS BEING THE M/2 00004170
C      LOW-ORDER BITS OF HALF1, AND THE M/2 LOW-ORDER BITS BEING 00004180
C      THE M/2 LOW-ORDER BITS OF HALF2.                  00004190
C      THIS IS DONE WITHOUT CAUSING AN INTEGER OVERFLOW AND 00004200
C      ASSUMING M IS EVEN, AND HIGH-ORDER (ISIZE-M)+M/2 BITS OF 00004210
C      EACH HALF ARE ZERO.                                00004220
      INTEGER HALF1,HALF2,ISUM(2),SHIFT1                 00004230
      COMMON/MACHIN/ISIZE,M,MAXPOS,MINNEG,MASK           00004240
      SHIFT1=2**((M/2)-1)                                00004250
C-----FORM JOIN=HALF2+(2**(M/2))*HALF1                 00004260
      IX=HALF1*SHIFT1                                     00004270
      CALL LADD(IX,IX,ISUM)                               00004280
      JOIN=ISUM(2)+HALF2                                  00004290
      RETURN                                               00004300
      END                                                  00004310

SUBROUTINE SPLIT(IX,HALF1,HALF2)                        00004320
CXXXXXXXXXXXXX ROUTINE TO CONVERT A RADIX 2**M DIGIT, IX TO TWO 00004330
C      RADIX 2**(M/2) DIGITS HALF1,HALF2. THAT IS THIS 00004340
C      ROUTINE SPLITS THE M-BIT NUMBER IX INTO HIGH-ORDER HALF 00004350
C      HALF1 AND LOW-ORDER HALF2 WITHOUT CAUSING AN INTEGER 00004360
C      OVERFLOW. THE HIGH-ORDER ISIZE-M BITS OF IX ARE ASSUMED 00004370
C      TO BE ZERO.                                        00004380
      INTEGER X,HALF1,HALF2,SHIFT,ISUM(2)               00004390
      COMMON/MACHIN/ISIZE,M,MAXPOS,MINNEG,MASK           00004400

```



```

SHIFT=2**(M/2)                                00004410
X=IX                                           00004420
IF(X.GE.0.AND.X.LE.MAXPOS) GO TO 2            00004430
C-----REMOVE SIGN BIT BEFORE DIVISION.      00004440
CALL LADD(MINNEG,X,ISUM)                      00004450
HALF1=ISUM(2)/SHIFT+SHIFT/2                  00004460
C-----FORM X MOD 2**(M/2)                   00004470
HALF2=MOD(ISUM(2),SHIFT)                     00004480
RETURN                                         00004490
2 HALF1=X/SHIFT                                00004500
HALF2=MOD(X,SHIFT)                            00004510
RETURN                                         00004520
END                                             00004530

SUBROUTINE LSUB(IU,IV,IW)                      00004540
CXXXXXXXXX SUBROUTINE TO PERFORM A LOGICAL    00004550
C          NUMBERS IU, IV. IW(2) IS THEIR     00004560
C          LOGICAL DIFFERENCE, IW(1) IS +1    00004570
C          IF THERE WAS A BORROW, ZERO        00004580
C          OTHERWISE. SEE SUBROUTINE LADD.     00004590
INTEGER IU,IV,IX,IY,IW(2),ISUM(2)           00004600
COMMON /MACHIN/ISIZE,M,MAXPOS,MINNEG,MASK    00004610
IX=IU                                         00004620
IY=IV                                         00004630
IW(1)=0                                       00004640
C-----SORT OUT SPECIAL CASES FOR LADD.      00004650
IF(IX.EQ.IY) GO TO 3                          00004660
IF(IY.EQ.0) GO TO 1                           00004670
IF(IY.EQ.MINNEG) GO TO 2                      00004680
CALL LADD(IX,MASK-IY,ISUM)                    00004690
IF(ISUM(1).EQ.0) IW(1)=1                     00004700
IW(2)=ISUM(2)                                00004710
RETURN                                         00004720
1 IW(2)=IX                                    00004730
RETURN                                         00004740
2 CALL LADD(IX,IY,ISUM)                       00004750
IF(ISUM(1).EQ.0) IW(1)=1                     00004760
IW(2)=ISUM(2)                                00004770
RETURN                                         00004780
3 IW(2)=0                                       00004790
RETURN                                         00004800
END                                             00004810

SUBROUTINE LADD(IU,IV,IW)                      00004820
CXXXXXXXXXXXXXXXXX SUBROUTINE TO PERFORM A    00004830
C          LOGICAL ADD OF THE LOW-ORDER      00004840
C          M BITS OF TWO COMPUTER WORDS, IU,  00004850
C          IV. IW(2) IS THEIR LOGICAL SUM,    00004860
C          IW(1) IS THE CARRY BIT VALUE. NO    00004870
C          INTEGER OVERFLOWS OCCUR EVEN IF    00004880
C          M=ISIZE. EXAMPLE - 0111 LADD 1011  00004890
C          YIELDS IW(1)=1, IW(2)=0010, FOR    00004900
C          M=4, MAXPOS=0111, MINNEG=1000,     00004910
C          MASK=0 OR 10000.                    00004920
C          INTEGER IX,IY,U,IV,IW(2),DIF,NEG,  00004930
C          POS,ISUM,BIG,SMALL,ISX,ISY,ISD     00004940
COMMON/MACHIN/ISIZE,M,MAXPOS,MINNEG,MASK     00004950
IX=IU                                         00004960
IY=IV                                         00004970
IW(1)=0                                       00004980
C-----ELIMINATE ZERO OPERANDS.              00004990
IF(IX.EQ.0.OR.IY.EQ.0) GO TO 2                00005000
C-----DETERMINE IF IX,IY IS ++, --, OR +-   00005010
ISX=ISIGNM(IX)                                00005020
ISY=ISIGNM(IY)                                00005030
IF(ISX.NE.ISY) GO TO 5                         00005040
IF(ISX.LT.0) GO TO 3                           00005050
C-----++ (BOTH OPERANDS HAVE HIGH-ORDER     00005060
C          BITS OFF)                          00005070
1 DIF=MIN0(IX,IY)-(MAXPOS-MAX0(IX,IY))        00005080
ISD=ISIGNM(DIF)                               00005090
IF(ISD.LE.0) IW(2)=IX+IY                      00005100
IF(ISD.GT.0) IW(2)=(MINNEG+DIF)-1             00005110
RETURN                                         00005120
C-----AT LEAST ONE OPERAND IS ZERO.          00005130
2 IW(2)=IX+IY                                00005140
RETURN                                         00005150
C----- -- (BOTH OPERANDS HAVE HIGH-ORDER    00005160
C          BITS ON)                           00005170
3 IW(1)=1                                       00005180
BIG=MAX0(IX,IY)                               00005190

```


ALGORITHM 537

Characteristic Values of Mathieu's Differential Equation [S22]

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Key Words and Phrases: Mathieu's differential equation, wave equation, characteristic values, eigenvalues, separation constants, Mathieu functions, ordinary Mathieu functions, modified Mathieu functions, elliptic cylinder functions, hyperbolic cylinder functions

CR Categories: 5.14, 5.17

Language: Fortran

DESCRIPTION

The algorithm calculates the characteristic values b of Mathieu's differential equation

$$d^2y/dx^2 + (b - S \cos^2 x)y = 0 \quad (1)$$

for solutions $y(x)$ of periodicity π or 2π . In this case b is a function of S and belongs to a countably infinite set of values for every S [2]. The knowledge of b is a prerequisite for calculating ordinary and modified Mathieu functions.

The program has been tested for real S -parameters between zero and 1000, and for a number of characteristic values of up to 24, depending on the value of S . (For negative S , simple equations exist which relate the corresponding b to those for positive S [2, p. xvii]). The calculated b are correct to at least 9 decimal places. Within the calculated finite subset of b , the characteristic values are arranged with increasing magnitude, the first value being the smallest of the infinite set. The algorithm uses a library subroutine to find the eigenvalues of a real symmetric tridiagonal matrix.

Four types of solutions $y(x)$ can be distinguished: An odd and an even solution with period 2π , and an odd and an even solution with period π . The corresponding characteristic values b are called bo_{2r+1} , be_{2r+1} , bo_{2r+2} ,¹ and be_{2r} . ($r = 0, 1, 2, \dots$; within each of the four sets, the smallest index is used for the smallest value of b .) Upon insertion of $y(x)$ in the form of Fourier series into eq. (1), a system of homogeneous, linear equations for the Fourier coefficients is obtained [2]. The determinant of the resulting square coefficient matrix of infinite order has to equal zero for nontrivial $y(x)$. This characteristic equation determining the

¹ Note the difference in the index notation for bo_{2r+2} compared with [2], where this set of characteristic values is called bo_{2r} . We have thus avoided unequal starting values for r among the four sets.

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characteristic values b may be written as

$$|\beta I - A| = 0, \tag{2}$$

where $\beta = b - S/2$ and I is the identity matrix. In the case of y being odd with a periodicity of 2π , the elements of the matrix A are found to be

$$a_{11} = 1 - S/4; a_{i-1,i} = a_{i,i-1} = S/4, a_{ii} = (2i - 1)^2 \text{ for } i \geq 2,$$

and all the other a_{ij} are equal to zero. Therefore, after adding $S/2$, the eigenvalues of the real symmetric tridiagonal matrix A (of unrestricted order) yield the characteristic values bo_{2r+1} .

Correspondingly, one can show that the be_{2r+1} , bo_{2r+2} , and be_{2r} are given by the eigenvalues of matrices which we call B , C and D , respectively. The nonzero elements of all four matrices have been compiled in Table I. When proceeding the way outlined above, the matrix D' determining the characteristic values be_{2r}

Table I. Nonzero Elements for Calculation of Characteristic Values

Characteristic value	Matrix elements
bo_{2r+1}	$a_{11} = 1 - S/4; a_{i-1,i} = a_{i,i-1} = S/4$ and $a_{ii} = (2i - 1)^2$ for $i \geq 2$
be_{2r+1}	$b_{11} = 1 + S/4; b_{i-1,i} = b_{i,i-1} = S/4$ and $b_{ii} = (2i - 1)^2$ and $i \geq 2$
bo_{2r+2}	$c_{ii} = 4i^2$ for $i \geq 1; c_{i,i-1} = c_{i-1,i} = S/4$ for $i \geq 2$
be_{2r}	$d_{12} = d_{21} = S/(8^{1/2}); d_{22} = 4$ $d_{ii} = 4(i - 1)^2$ and $d_{i-1,i} = d_{i,i-1} = S/4$ for $i \geq 3$

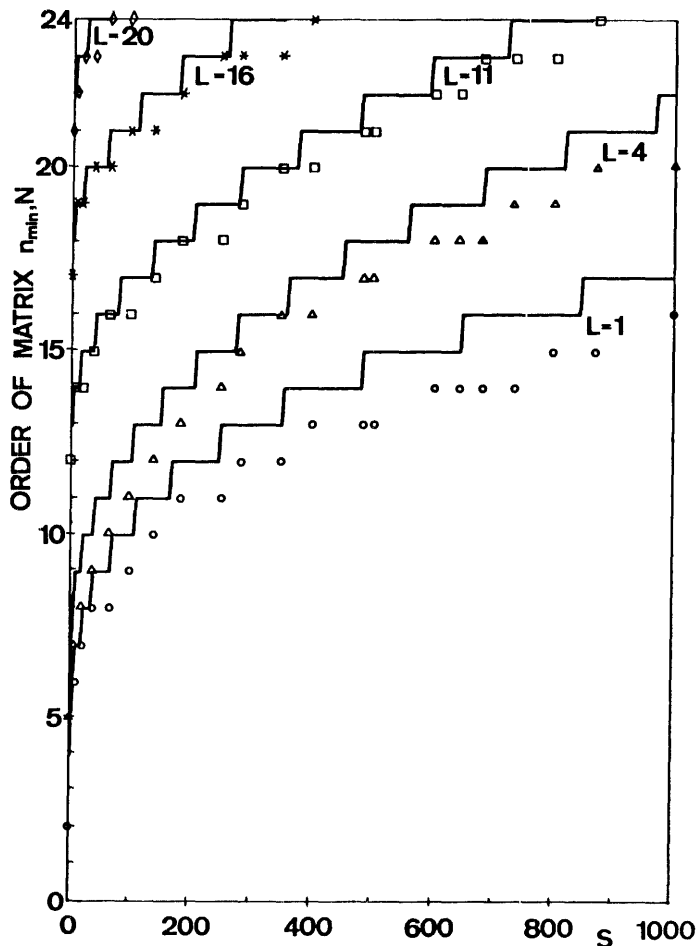


Fig. 1. Order of matrix (A , B , C , or D) necessary to give the first L characteristic values b correct to nine decimal places. Data points (\circ , Δ , \square , $*$, and \diamond) represent numerically found n_{\min} in case of matrix D , while the full lines represent the function $N(S, L)$ according eq. (3). For clarity, only five values of L are presented.

turns out to be tridiagonal but unsymmetric. The only reason for unsymmetry is that $d'_{12} = S/4$ is unequal $d'_{21} = S/2$. Proper similarity transformation produces a symmetric tridiagonal matrix D with $d_{12} = d_{21} = S/(8^{1/2})$ and all other elements remain unchanged. This allows the use of one and the same library subroutine for determining the eigenvalues of all four matrices.

Next, one has to establish the minimum order n_{\min} of the matrix A (or B, C, D) such as to yield a number L ($L < n_{\min}$) of characteristic values with given accuracy. In a modification of the program presented here, and using a trial and error method, we determined n_{\min} as a function of both S and L for a b -accuracy of 9 decimal places. This was done for some discrete values of S and for integer L in the domain $0 \leq S \leq 1000$ and $0 \leq n_{\min} \leq 24$. The results are shown in Fig. 1 for some selected values of L . Then an analytic function $N(S, L)$ was developed which closely approximates n_{\min} , but ensures $N \geq n_{\min}$:

$$N(S, L) = \text{int}((0.17 + 2.1 \exp(-0.24 L)) S^{(0.77-5/(9.5+L))} + L + 2.8). \quad (3)$$

The function $N(S, L)$ is presented in Fig. 1 as full line for the same parameters L as the numerically evaluated $n_{\min}(S, L)$. There are only minor differences in n_{\min} required for A, B, C , and D , and eq. (3) has been determined as to fulfill $N \geq n_{\min}$ for all four matrices; therefore, it will be used for all four sets of characteristic values.

EXAMPLES

Find $bo_1, bo_3, \dots, bo_{29}$ for $S = 2$, also be_1, be_3, \dots, be_7 for $S = 1000$; in addition, $bo_2, bo_4, \dots, bo_{38}$ for $S = 100$, and $be_0, be_2, \dots, be_{18}$ for $S = 0.1$. The results computed with a CDC CYBER/16 are given in Table II. Part of these results can be compared with tabulated values [2] and show agreement.

ACKNOWLEDGMENT

Thanks are due to the Computer Center of the Technische Universität Wien, which provided free computer time, and to R. Welser for his assistance.

Table II

S =	2.00	CHARACTERISTIC VALUES	S = 1000.00	CHARACTERISTIC VALUES
		BO 1 = 1.470654355		BE 1 = 93.599796781
		BO 3 = 10.013719839		BE 3 = 214.914687921
		BO 5 = 26.005209010		BE 5 = 331.837262869
		BO 7 = 50.002604266		BE 7 = 444.079857151
		BO 9 = 82.001562520		
		BO11 = 122.001041673		
		BO13 = 170.000744050		
		BO15 = 226.000558037		
		BO17 = 290.000434028		
		BO19 = 362.000347222		
		BO21 = 442.000284091		
		BO23 = 530.000236742		
		BO25 = 626.000200321		
		BO27 = 730.000171703		
		BO29 = 842.000148899		
			S = 100.00	CHARACTERISTIC VALUES
				BO 2 = 28.685139778
				BO 4 = 62.986489953
				BO 6 = 91.801071792
				BO 8 = 119.057988351
				BO10 = 153.225680742
				BO12 = 196.207674647
				BO14 = 247.611164916
				BO16 = 307.229284862
				BO18 = 374.969344509
				BO20 = 450.784185569
				BO22 = 534.647547063
				BO24 = 626.547802911
				BO26 = 726.463163268
				BO28 = 834.399234146
				BO30 = 950.347693024
				BO32 = 1074.305531439
				BO34 = 1206.270602586
				BO36 = 1346.241340972
				BO38 = 1494.216583111
S =	.10	CHARACTERISTIC VALUES		
		BE 0 = .049687521		
		BE 2 = 4.050260395		
		BE 4 = 16.050020834		
		BE 6 = 36.050008929		
		BE 8 = 64.050004960		
		BE10 = 100.050003157		
		BE12 = 144.050002185		
		BE14 = 196.050001603		
		BE16 = 256.050001225		
		BE18 = 324.050000967		

REFERENCES

[Note. Reference [1] is not mentioned in the text.]

1. Goos, G., AND HARTMANIS, J. Eds., *Lecture Notes in Computer Science, vol. 6*, Springer, New York, 1976.
2. National Bureau of Standards, *Tables Relating to Mathieu Functions*. Cambridge U. Press, New York, 1951.

ALGORITHM

```

SUBROUTINE CHARMA(KIND, S, L, NMAX, VAL, SUD, IA)          00000010
C SUBROUTINE CHARMA CALCULATES THE CHARACTERISTIC VALUES OF 00000020
C MATHIEU-S DIFFERENTIAL EQUATION FOR ODD OR EVEN SOLUTIONS 00000030
C WITH PERIODICITY PI OR 2*PI.                          00000040
C A LIBRARY SUBROUTINE HAS TO BE ATTACHED TO CALCULATE THE 00000050
C EIGENVALUES OF A REAL SYMMETRIC TRIDIAGONAL MATRIX. IN THE 00000060
C PRESENTED FORM, THIS IS THE EISPACK ROUTINE IMTQL1.     00000070
C INPUT..                                                00000080
C   KIND AN INTEGER CHARACTERIZING THE KIND OF CH. V.    00000090
C   IF KIND=1, THE CH. V. BO 1, BO 3,... FOR ODD SOLUTIONS 00000100
C   WITH PERIODICITY 2*PI ARE CALCULATED                 00000110
C   IF KIND=2, THE CH. V. BE 1, BE 3,... FOR EVEN SOLUTIONS 00000120
C   WITH PERIODICITY 2*PI ARE CALCULATED                 00000130
C   IF KIND=3, THE CH. V. BO 2, BO 4,... FOR ODD SOLUTIONS 00000140
C   WITH PERIODICITY PI ARE CALCULATED                   00000150
C   IF KIND=4, THE CH. V. BE 0, BE 2,... FOR EVEN SOLUTIONS 00000160
C   WITH PERIODICITY PI ARE CALCULATED                   00000170
C   CHARMA DOES NOT DESTROY KIND                         00000180
C   S REAL NON-NEGATIVE VARIABLE, THE PARAMETER OF THE 00000190
C   DIFFERENTIAL EQUATION. FOR S UP TO 1000 AN ACCURACY OF 00000200
C   9 DECIMAL PLACES IS GUARANTEED FOR THE CHARACTERISTIC 00000210
C   VALUES IF THE NUMBER OF THE CH. V. IS NOT TOO HIGH. 00000220
C   (SEE FIG. 1 OF DESCRIPTION)                          00000230
C   L INTEGER VARIABLE, THE NUMBER OF CHARACTERISTIC VALUES 00000240
C   TO BE CALCULATED. IT CAN BE NO LARGER THAN NMAX.    00000250
C   NMAX INTEGER VARIABLE, MAXIMUM DIMENSION OF THE MATRIX USED 00000260
C   FOR THE CALCULATION. TO MAKE USE OF THE FULL TESTED 00000270
C   DOMAIN WITH AN ACCURACY OF 9 DECIMAL PLACES, IT SHOULD 00000280
C   BE AT LEAST 24.                                     00000290
C OUTPUT..                                              00000300
C   VAL REAL ONE-DIMENSIONAL ARRAY OF DIMENSION (NMAX). 00000310
C   INITIALLY IT CONTAINS THE DIAGONAL ELEMENTS OF THE 00000320
C   COEFFICIENT MATRIX. ON EXIT, ITS FIRST L ELEMENTS WILL 00000330
C   CONTAIN THE CH. V.                                   00000340
C   IA INTEGER VARIABLE USED AS A FAILURE INDICATOR. IF, ON 00000350
C   EXIT, IA=0, NO FAILURE WAS DETECTED. IF IA=1, S WAS 00000360
C   NEGATIVE. IF IA=2, L WAS CHOSEN TOO BIG, REQUIRING A 00000370
C   LARGER NMAX. IF IA=3, THE LIBRARY SUBROUTINE IMTQL1 00000380
C   DID NOT FIND ALL EIGENVALUES. FOR IA=1, 2, OR 3, NO 00000390
C   CHARACTERISTIC VALUES WERE CALCULATED. IF IA=4, 00000400
C   S .GT. 1000. IF IA=5, L WAS TOO LARGE (FOR THE GIVEN 00000410
C   S), REQUIRING AN ORDER N OF THE COEFFICIENT MATRIX 00000420
C   WHICH EXCEEDS THE TESTED DOMAIN (N .LE. 24). FOR IA=4 00000430
C   OR 5 THE CALCULATION WAS EXECUTED, BUT ACCURACY OF THE 00000440
C   CHARACTERISTIC VALUES TO NINE DECIMAL PLACES IS NOT 00000450
C   GUARANTEED. IN THESE CASES, ACCURACY MAY BE CHECKED AS 00000460
C   FOLLOWS. IF L=L1 PRODUCES IA=4 OR 5, TAKE A VALUE L=L2, 00000470
C   L2=L1+1, COMPUTE VAL (IA WILL AGAIN BE 4 OR 5) AND COMPUTE 00000480
C   THE DIFFERENCE OF THE FIRST L1 MEMBERS OF BOTH SEQUENCES. 00000490
C   IF THE ERROR TEST IS PASSED, ACCEPT THE ANSWER, IF NOT TAKE 00000500
C   A VALUE L=L3=L2+1, ETC...                            00000510
C OTHER PARAMETERS..                                    00000520
C   SUD A ONE-DIMENSIONAL REAL ARRAY OF DIMENSION (NMAX), 00000530
C   INITIALLY CONTAINING IN ITS POSITIONS (2), (3),..., (N) 00000540
C   THE SUBDIAGONAL ELEMENTS OF THE COEFFICIENT MATRIX . 00000550
C   N INTEGER VARIABLE, THE ORDER OF THE COEFFICIENT MATRIX 00000560
C   IB AN INTEGER VARIABLE TO TEST THE SUCCESS OF IMTQL1. IF, 00000570
C   ON EXIT, IB=0, IMTQL1 HAS DETERMINED ALL EIGENVALUES 00000580
C   WITHIN 30 ITERATIONS.                               00000590
C   REAL FL, S, SUD, VAL                                00000600
C   INTEGER I, IA, IB, IOUT, KIND, L, N, NMAX           00000610
C   DIMENSION VAL(NMAX), SUD(NMAX)                    00000620
C   IA = 0                                             00000630
C TEST FOR NEGATIVE S                                  00000640
C   IF (S.GE.0.0E0) GO TO 10                          00000650

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IA = 1                                00000660
RETURN                                00000670
C TEST FOR S GREATER 1000. IF TRUE, CALCULATION IS CONTINUED 00000680
C BUT ACCURACY OF THE CHARACTERSTIC VALUES TO NINE DECIMAL 00000690
C PLACES CANNOT BE GUARANTEED. SEE COMMENT ON IA IN OUTPUT LIST 00000700
C FOR IMPROVEMENT.                   00000710
10 IF (S.LE.1000.0E0) GO TO 20        00000720
IA = 4                                00000730
C DETERMINE NECESSARY ORDER OF MATRIX TO ACHIEVE AN ACCURACY 00000740
C OF 9 DECIMAL PLACES FOR THE CHARACTERISTIC VALUES.        00000750
20 FL = FLOAT(L)                       00000760
N = INT((0.17E0+2.1E0*EXP(-0.24E0*FL))*S**((0.77E0-5.0E0/ 00000770
* (9.5E0+FL))+FL+2.8E0))              00000780
C TEST FOR SUFFICIENT LARGE NMAX      00000790
IF (N.LE.NMAX) GO TO 30               00000800
IA = 2                                00000810
RETURN                                00000820
C TEST WHETHER N IS WITHIN TESTED DOMAIN FOR WHICH ACCURACY TO 9 00000830
C DECIMAL PLACES IS GUARANTEED. IF NOT, CALCULATION IS CONTINUED 00000840
C BUT ACCURACY OF THE CHARACTERSTIC VALUES TO NINE DECIMAL 00000850
C PLACES CANNOT BE GUARANTEED. SEE COMMENT ON IA IN OUTPUT LIST 00000860
C FOR IMPROVEMENT.                   00000870
30 IF (N.LE.24) GO TO 40              00000880
IA = 5                                00000890
C BRANCH ACCORDING TO DESIRED SOLUTION: 00000900
C IF KIND=1, USE MATRIX CALLED A IN THE DESCRIPTION 00000910
C IF KIND=2, USE MATRIX CALLED B IN THE DESCRIPTION 00000920
C IF KIND=3, USE MATRIX CALLED C IN THE DESCRIPTION 00000930
C IF KIND=4, USE MATRIX CALLED D IN THE DESCRIPTION 00000940
40 GO TO (50, 60, 90, 110), KIND      00000950
C STORE DIAGONAL ELEMENTS OF THE COEFFICIENT MATRIX IN VAL AND 00000960
C SUBDIAGONAL ELEMENTS IN SUD(2), SUD(3),..., SUD(N).        00000970
50 VAL(1) = 1.0E0 - S/4.0E0           00000980
GO TO 70                               00000990
60 VAL(1) = 1.0E0 + S/4.0E0           00010000
70 DO 80 I=2,N                         00010010
VAL(I) = FLOAT((2*I-1)**2)            00010020
SUD(I) = S/4.0E0                     00010030
80 CONTINUE                            00010040
GO TO 130                              00010050
90 VAL(1) = 4.0E0                      00010060
DO 100 I=2,N                           00010070
VAL(I) = FLOAT((2*I)**2)              00010080
SUD(I) = S/4.0E0                     00010090
100 CONTINUE                           00010100
GO TO 130                              00010110
110 VAL(1) = 0.0E0                      00010120
VAL(2) = 4.0E0                         00010130
SUD(2) = S/SQRT(8.0E0)                 00010140
DO 120 I=3,N                           00010150
VAL(I) = FLOAT((2*(I-1))**2)          00010160
SUD(I) = S/4.0E0                     00010170
120 CONTINUE                           00010180
130 CALL IMTQL1(N, VAL, SUD, IB)        00010190
C TEST FOR SUCCESSFUL IMTQL1           00010200
IF (IB.EQ.0) GO TO 140                 00010210
IA = 3                                  00010220
RETURN                                  00010230
C ADD S/2 TO THE EIGENVALUES OF THE MATRIX TO GET THE CH. V. 00010240
140 DO 150 I=1,L                       00010250
VAL(I) = VAL(I) + S/2.0E0             00010260
150 CONTINUE                           00010270
RETURN                                  00010280
END                                     00010290

C PROGRAM TECHV (INPUT,OUTPUT,TAPE5=INPUT,TAPE6=OUTPUT)      00010300
C THIS IS A PROGRAM TO TEST THE SUBROUTINE CHARMA WHICH CALCULATES THE 00010310
C CHARACTERISTIC VALUES OF MATHIEU'S DIFFERENTIAL EQUATION FOR ODD OR 00010320
C EVEN SOLUTIONS WITH PERIODICITY PI OR 2*PI.                 00010330
C INPUT..                                                       00010340
C S A REAL VARIABLE, THE PARAMETER OF THE DIFFERENTIAL EQUATION 00010350
C L AN INTEGER, THE NUMBER OF CHARACTERISTIC VALUES TO BE    00010360
C CALCULATED                                                    00010370
C OUTPUT..                                                       00010380
C IA AN INTEGER USED AS A FAILURE INDICATOR. IF, ON EXIT, IA=0, NO 00010390
C FAILURE WAS DETECTED. IF IA=1 OR IA=2 OR IA=3, A FAILURE WAS 00010400

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C REGISTERED AND NO CHARACTERISTIC VALUES WERE CALCULATED. 00001410
C VAL A ONE-DIMENSIONAL ARRAY USED BY CHARMA AND BY A LIBRARY 00001420
C SUBROUTINE. THE DIMENSION OF VAL SHOULD BE EQUAL OR GREATER 00001430
C THAN 24 TO ACHIEVE FULL ACCURACY OVER THE WHOLE TESTED DOMAINE 00001440
C OF CHARMA. ON EXIT, VAL CONTAINS THE CHARACTERISTIC VALUES. 00001450
C INDE INTEGER VARIABLE, THE INDEX OF THE CHARACTERISTIC VALUES 00001460
C KIND AN INTEGER CHARACTERIZING THE KIND OF CHARACTERISTIC VALUE 00001470
C TO BE CALCULATED 00001480
C SUD A ONE-DIMENSIONAL ARRAY USED BY CHARMA AND BY A LIBRARY 00001490
C SUBROUTINE. THE DIMENSION OF SUD SHOULD BE EQUAL OR GREATER 00001500
C THAN 24 TO ACHIEVE FULL ACCURACY OVER THE WHOLE TESTED DOMAINE 00001510
C OF CHARMA. 00001520
C THE FOLLOWING DATA CARDS HAVE BEEN USED FOR THIS TEST PROGRAM (S=-1. 00001530
C SIGNALS END OF DATA) 00001540
C 2. 15 00001550
C1001. 4 00001560
C 0.01 25 00001570
C 10. 27 00001580
C -6. 3 00001590
C -1. 00001600
      DIMENSION VAL(28), SUD(28) 00001610
      10 READ (5,99999) S, L 00001620
      IF (S.EQ.-1.) GO TO 80 00001630
C TEST SUBROUTINE FOR ALL FOUR KIND OF SOLUTIONS 00001640
      DO 70 KIND=1,4 00001650
      CALL CHARMA(KIND, S, L, 28, VAL, SUD, IA) 00001660
      IF (IA.GT.0) WRITE (6,99993) S, L, IA 00001670
      IF ((IA.GT.0) .AND. (IA.LT.4)) GO TO 10 00001680
      WRITE (6,99998) S 00001690
C PRINT THE FIRST L ELEMENTS OF VAL, WHICH ARE THE CHARACTERISTIC VALUES 00001700
      DO 60 I=1,L 00001710
C GIVE NAMES AND INDICES OF CH. V. ACCORDING KIND OF SOLUTION 00001720
      GO TO (20, 30, 40, 50), KIND 00001730
      20 INDE = 2*I - 1 00001740
      WRITE (6,99997) INDE, VAL(I) 00001750
      GO TO 60 00001760
      30 INDE = 2*I - 1 00001770
      WRITE (6,99996) INDE, VAL(I) 00001780
      GO TO 60 00001790
      40 INDE = 2*I 00001800
      WRITE (6,99995) INDE, VAL(I) 00001810
      GO TO 60 00001820
      50 INDE = 2*I - 2 00001830
      WRITE (6,99994) INDE, VAL(I) 00001840
      60 CONTINUE 00001850
      70 CONTINUE 00001860
      GO TO 10 00001870
      80 CONTINUE 00001880
      STOP 00001890
99999 FORMAT (F8.2, 2X, I2) 00001900
99998 FORMAT (1H0, 17X, 3HS =, F8.2, 2X, 19HCHARACTERISTIC VALU, 00001910
* 2HES/) 00001920
99997 FORMAT (31X, 2HBO, I2, 2H =, F14.9) 00001930
99996 FORMAT (31X, 2HBE, I2, 2H =, F14.9) 00001940
99995 FORMAT (31X, 2HBO, I2, 2H =, F14.9) 00001950
99994 FORMAT (31X, 2HBE, I2, 2H =, F14.9) 00001960
99993 FORMAT (1H0, 30X, 10HDATA - S =, F8.2/38X, 3HL =, 00001970
* I3//26X, 15HERROR FLAG IA =, I3//) 00001980
      END 00001990
C 00002000
C ----- 00002010
C 00002020
C SUBROUTINE IMTQL1(N,D,E,IERR) 00002030
C 00002040
C INTEGER I,J,L,M,N,II,MML,IERR 00002050
C REAL D(N),E(N) 00002060
C REAL B,C,F,G,P,R,S,MACHEP 00002070
C REAL SQRT,ABS,SIGN 00002080
C 00002090
C THIS SUBROUTINE IS A TRANSLATION OF THE ALGOL PROCEDURE IMTQL1, 00002100
C NUM. MATH. 12, 377-383(1968) BY MARTIN AND WILKINSON, 00002110
C AS MODIFIED IN NUM. MATH. 15, 450(1970) BY DUBRULLE. 00002120
C HANDBOOK FOR AUTO. COMP., VOL.II-LINEAR ALGEBRA, 241-248(1971). 00002130
C 00002140
C THIS SUBROUTINE FINDS THE EIGENVALUES OF A SYMMETRIC 00002150
C TRIDIAGONAL MATRIX BY THE IMPLICIT QL METHOD. 00002160

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C          00002170
C      ON INPUT-          00002180
C          00002190
C      N IS THE ORDER OF THE MATRIX,          00002200
C          00002210
C      D CONTAINS THE DIAGONAL ELEMENTS OF THE INPUT MATRIX,          00002220
C          00002230
C      E CONTAINS THE SUBDIAGONAL ELEMENTS OF THE INPUT MATRIX          00002240
C      IN ITS LAST N-1 POSITIONS. E(1) IS ARBITRARY.          00002250
C          00002260
C      ON OUTPUT-          00002270
C          00002280
C      D CONTAINS THE EIGENVALUES IN ASCENDING ORDER. IF AN          00002290
C      ERROR EXIT IS MADE, THE EIGENVALUES ARE CORRECT AND          00002300
C      ORDERED FOR INDICES 1,2,...IERR-1, BUT MAY NOT BE          00002310
C      THE SMALLEST EIGENVALUES,          00002320
C          00002330
C      E HAS BEEN DESTROYED,          00002340
C          00002350
C      IERR IS SET TO          00002360
C      ZERO          FOR NORMAL RETURN,          00002370
C      J          IF THE J-TH EIGENVALUE HAS NOT BEEN          00002380
C                  DETERMINED AFTER 30 ITERATIONS.          00002390
C          00002400
C      QUESTIONS AND COMMENTS SHOULD BE DIRECTED TO B. S. GARBOW,          00002410
C      APPLIED MATHEMATICS DIVISION, ARGONNE NATIONAL LABORATORY          00002420
C          00002430
C      -----          00002440
C          00002450
C      ***** MACHEP IS A MACHINE DEPENDENT PARAMETER SPECIFYING          00002460
C      THE RELATIVE PRECISION OF FLOATING POINT ARITHMETIC.          00002470
C          00002480
C          *****          00002490
C      MACHEP = 2.**(-26)          00002500
C          00002510
C      IERR = 0          00002520
C      IF (N .EQ. 1) GO TO 1001          00002530
C          00002540
C      DO 100 I = 2, N          00002550
1000 E(I-1) = E(I)          00002560
C          00002570
C      E(N) = 0.0          00002580
C          00002590
C      DO 290 L = 1, N          00002600
C          J = 0          00002610
C      ***** LOOK FOR SMALL SUB-DIAGONAL ELEMENT *****          00002620
105  DO 110 M = L, N          00002630
C          IF (M .EQ. N) GO TO 120          00002640
C          IF (ABS(E(M)) .LE. MACHEP * (ABS(D(M)) + ABS(D(M+1))))          00002650
C      X          GO TO 120          00002660
110  CONTINUE          00002670
C          00002680
C      120  P = D(L)          00002690
C          IF (M .EQ. L) GO TO 215          00002700
C          IF (J .EQ. 30) GO TO 1000          00002710
C          J = J + 1          00002720
C          ***** FORM SHIFT *****          00002730
C          G = (D(L+1) - P) / (2.0 * E(L))          00002740
C          R = SQRT(G*G+1.0)          00002750
C          G = D(M) - P + E(L) / (G + SIGN(R,G))          00002760
C          S = 1.0          00002770
C          C = 1.0          00002780
C          P = 0.0          00002790
C          MML = M - L          00002800
C          ***** FOR I=M-1 STEP -1 UNTIL L DO -- *****          00002810
C      DO 200 II = 1, MML          00002820
C          I = M - II          00002830
C          F = S * E(I)          00002840
C          B = C * E(I)          00002850
C          IF (ABS(F) .LT. ABS(G)) GO TO 150          00002860
C          C = G / F          00002870
C          R = SQRT(C*C+1.0)          00002880
C          E(I+1) = F * R          00002890
C          S = 1.0 / R          00002900
C          C = C * S          00002910
C          GO TO 160          00002920

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```

150      S = F / G
          R = SQRT(S*S+1.0)
          E(I+1) = G * R
          C = 1.0 / R
          S = S * C
160      G = D(I+1) - P
          R = (D(I) - G) * S + 2.0 * C * B
          P = S * R
          D(I+1) = G + P
          G = C * R - B
200      CONTINUE
C
          D(L) = D(L) - P
          E(L) = G
          E(M) = 0.0
          GO TO 105
C ***** ORDER EIGENVALUES *****
215      IF (L .EQ. 1) GO TO 250
C ***** FOR I=L STEP -1 UNTIL 2 DO -- *****
          DO 230 II = 2, L
              I = L + 2 - II
              IF (P .GE. D(I-1)) GO TO 270
              D(I) = D(I-1)
230      CONTINUE
C
250      I = 1
270      D(I) = P
290      CONTINUE
C
          GO TO 1001
C ***** SET ERROR -- NO CONVERGENCE TO AN
C ***** EIGENVALUE AFTER 30 ITERATIONS *****
1000      IERR = L
1001      RETURN
C ***** LAST CARD OF IMTQL1 *****
          END

```

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00002930
00002940
00002950
00002960
00002970
00002980
00002990
00003000
00003010
00003020
00003030
00003040
00003050
00003060
00003070
00003080
00003090
00003100
00003110
00003120
00003130
00003140
00003150
00003160
00003170
00003180
00003190
00003200
00003210
00003220
00003230
00003240
00003250
00003260
00003270
00003280

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ALGORITHM 538

Eigenvectors and Eigenvalues of Real Generalized Symmetric Matrices by Simultaneous Iteration [F2]

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Key Words and Phrases: eigenvalue, eigenvector, sparse matrix, diagonalizable matrix, simultaneous iteration, Fortran program

CR Categories: 4.6, 5.14

Language: Fortran

DESCRIPTION

The program presented here is an implementation of the simultaneous iteration algorithm [2] for calculating the eigenvalues largest in magnitude and corresponding eigenvectors of a real matrix symmetric relative to a prescribed inner product. Let $\text{ip}(n, w, z)$ denote an inner product in the space of real column n -tuples and let the real n -square matrix C satisfy $\text{ip}(n, Cw, z) = \text{ip}(n, w, Cz)$. Then C is *symmetric relative to ip*, and if the n -square positive definite matrix B satisfies $\text{ip}(n, w, z) = w^T Bz$ then C is *B-symmetric*. The equation $BC = C^T B$ characterizes the *B-symmetry* of C . Given an optional set of p initial approximate eigenvectors of a real n -square *B-symmetric* matrix C corresponding to p eigenvalues of C largest in magnitude, the program calculates em eigenvalues and em corresponding eigenvectors, $0 \leq em < p \leq n$, to a precision dependent on the structure of C and on a prescribed tolerance eps . The matrix B is presented to the program as an independently prepared real function subprogram which calculates $\text{ip}(n, w, z) = w^T Bz$ given column n -vectors w and z . The matrix C is presented as an independently prepared subroutine subprogram $\text{op}(n, z, w)$ which when given an n -vector z computes its image $w = Cz$. The program is an outgrowth of a literal Fortran translation [5] of the Algol procedure *ritzit* [8] to which it is substantially equivalent when $C = C^T$ and $\text{ip}(n, w, z) = w^T z$, the standard inner product. But depending on the choice of B and C , the present program enables the direct treatment of a wide variety of symmetric eigenproblems.

Let $A = A^T$ and $B = B^T$ denote n -square real matrices and let σ be real. If B is positive definite then the matrix $C = B^{-1}(A - \sigma B)$ is *B-symmetric*, and the program computes eigenvalues farthest from σ of the eigenproblem $Au = \lambda Bu$ and corresponding eigenvectors. Implementation of $\text{op}(n, z, w)$ here consists in providing for the appropriate solution for w of the linear system $Bw = (A - \sigma B)z$. Alternatively, selection of op to solve the system $(A - \sigma B)w = Bz$ for w enables

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the calculation by simultaneous inverse iteration of the eigenvalues nearest to σ and their eigenvectors. Implications for large sparse systems for which the Cholesky factorization [6] of B is impractical are clear. The user may wish to supplement the following outline of the operation of the program by consulting the description of the Algol procedure *ritzit* in [8] or [10] as well as a review of the mathematical foundations of simultaneous iteration in [4] and [7].

Let the eigenvalues $d_1, \dots, d_p, d_{p+1}, \dots, d_n$ of C be arranged in order of descending absolute value and let E_p denote the direct sum of the distinct eigenspaces corresponding to d_1, \dots, d_p . Let X_0 denote an n -by- p matrix having a p -dimensional column space not orthogonal relative to ip to any eigenvector in E_p . Simultaneous iteration is based on the observation that if $|d_p| > |d_{p+1}|$, the columns of the matrix $X_{k+m} = C^m X_k$ tend to a basis of E_p as $ks = k + m$ increases. But in practice all of the columns of X_{ks} tend toward the eigenspace E_1 causing loss of information concerning the residual eigenvectors. To counter this tendency, set

$$X_{k+m} = C^m X_k R_{k+m}^{-1} \quad (1)$$

where the p -square upper triangular matrix R_{ks} is constructed together with X_{ks} by the Gram-Schmidt process to render the columns of X_{ks} orthonormal relative to ip. Now the i th column vector of X_{ks} converges to the i th eigenvector of C at a rate proportional to $\max_{2 \leq i \leq p} (|d_i/d_{i-1}|, |d_{i+1}/d_i|)$. Clearly this convergence will be delayed in the presence of eigenvalue clustering. But if $|d_p|/|d_{p+1}|$ is not too small, the column space of X_{ks} will contain a good approximation to the i th eigenvector even when ks is small.

In order to recover this approximation, a modified Rayleigh-Ritz process is employed. Let Q_{ks} denote an orthogonal matrix which diagonalizes the p -square symmetric matrix $R_{ks} R_{ks}^T$. Then the i th column vector of

$$X_{k+1} = C X_k R_{k+1}^{-1} Q_{k+1} \quad (2)$$

converges to the i th eigenvector of C at a rate proportional to $|d_{p+1}/d_i|$ while the entries of the diagonal matrix computed with Q_{ks} and properly ordered offer close approximations to d_1^2, \dots, d_p^2 . The true signed eigenvalues need only be computed at termination by diagonalizing the leading $(p-1)$ -square principal submatrix of $X_{ks}^T B C X_{ks}$, the eigenproblem for C projected on E_{p-1} relative to ip.

The program determines a strategy for employing the devices (1) and (2) based on the distribution of the leading p eigenvalues of C upon which the convergence rate ultimately depends. The selection of values m in eq. (1) is particularly important in this regard in that $C^m X_k$ is replaced by the m th Chebychev polynomial on the interval $[-e, e]$ evaluated by a special three-term recurrence relation and permitting accelerated convergence when values of m are continually large; e is the current value of d_p . As a result the convergence quotient lies between $|d_p/d_{em}|$ and $\exp(-\text{arc cosh } |d_{em}/d_p|)$. It is nearer to the first value if $|d_1/d_{em}|$ is large and nearer to the second if the latter quotient is close to 1.

As the iteration proceeds through a maximum of $|km|$ iteration steps (km is a program parameter), acceptance tests for the eigenvalues and eigenvectors are conducted following each of the Rayleigh-Ritz steps (2). As soon as the relative increase of $|d_{h+1}|$ is smaller than $\text{eps}/10$, then d_{h+1} is accepted and h , the number of previously accepted eigenvalues, is increased by 1. Eigenvectors are accepted in groups of one or more corresponding to clusters of accepted eigenvalues nearly equal in magnitude. If g eigenvectors have already been accepted, let d_{g+1}, \dots, d_l denote such a cluster. For all $j, g+1 \leq j \leq l$, denote by y_j the projection relative to ip of the image Cx_j of the t th column x_j of X_{ks} on the linear closure of x_1, \dots, x_l . Set $f_i = \max_j \|Cx_j - y_j\| / \|Cx_j\|$ for $i = g+1, \dots, l$ where the indicated norm is the Euclidean norm or 2-norm relative to ip. If $|d_i| f_i / (|d_i| - e)$ is smaller than eps then all the $x_j, j = g+1, \dots, l$, are accepted as eigenvectors and g is increased to l . The error quantities f_i are systematically discounted in accordance with the convergence properties of the algorithm to permit convergence in the presence of excessive roundoff error or in case the parameter eps is prescribed unrealistically

small. Having determined g eigenvectors, the iteration continues with $p - g$ remaining columns of X_{ks} until either em eigenvectors have been calculated or $|km|$ has been exceeded. The program may reduce em if it detects either no progress in convergence of eigenvectors corresponding to smaller eigenvalues or lack of stability in the behavior of larger eigenvalues.

We list the principal differences between *ritzit* and the program given here.

(1) The procedure *inprod* for calculating standard inner products was removed and the procedure *ip* was introduced where appropriate.

(2) The procedure *jacobi* for calculating the solutions of the eigenproblem for $(p - g)$ -square and $(p - 1)$ -square symmetric matrices was replaced by calls to the EISPACK [9] subroutines TRED2 and IMTQL2, primarily to save space.

(3) The procedure *random* for calculating random column n -vectors of the matrix X_{ks} was replaced by in-line code which references a Fortran function RANF. RANF returns uniformly distributed random REAL values from the interval $(0, 1)$, one per function reference, given any one argument of any type. It is provided by the user.

(4) The procedure *orthog* to perform Gram-Schmidt orthogonalization of the columns of X_{ks} was replaced by internally linked in-line code. In attempting to control potential underflow within *orthog* in a machine independent fashion, *ritzit* calculates the machine precision mc but assumes in usage that out-of-range values underflow gracefully to zero, a machine dependent characteristic. The present program utilizes a single REAL machine dependent constant MT, the ratio of the smallest Fortran representable positive value to the machine precision, to test for this condition and upon its detection to take appropriate measures.

(5) In its Algol implementation *ritzit* requires approximately $(p + 3)n + 2p^2 + 5p$ storage locations in excess of those required by the program. Economies resulting in part from (2) above [3] have reduced this requirement to $(p + 2)n + p^2 + 4p$ in the program presented here. All working storage is confined to a single array of $2n + p^2 + 3p$ locations.

(6) The value of km as an input parameter, set to $|km|$ during program execution, is finally replaced by the value of ks as an output parameter, the number of iteration steps used in the calculation of em eigenvectors.

(7) The program given here retains unchanged the reference to a user supplied procedure *inf* as a window on program execution. However, the one variable involving eps is periodically redefined to enable effect control of eps from *inf* or from *ip* or *op* should this prove desirable.

The testing procedures developed for the present Fortran program parallel its evolution from a research tool, which conformed closely to its Algol parent, to present form. Early testing was concentrated on duplicating the tests furnished with *ritzit* and in eliminating errors in interpretation and translation of the Algol code. This was done for the most part on the CDC 6600 using Fortran Extended, Version 3, under the SCOPE 3.3 operating system. Upon completion of this first phase, the resulting program was distributed as SUBROUTINE RITZIT with a locally developed library of Fortran linear algebra routines [5]. This same program served as a basis of *ritzit* translations for IBM 360/370 processors [1, 3] whose preparation uncovered several bugs in the RITZIT code and suggested worthwhile modifications. A second phase of testing involved the development of a package of auxiliary Fortran programs for use with SUBROUTINE RITZIT to solve the eigenvalue problem $Au = \lambda Bu$ through methods depending on Cholesky factorization where A and B may either be full matrices or sparse and banded. This phase was conducted on the CDC 6600 using Fortran Extended, Version 4, under SCOPE 3.4.

Systematic testing of the present program, SUBROUTINE SIMITZ, has been accomplished in part with the aid of driver program TESTB which generates a symmetric band matrix A and a lower triangular band matrix T of prescribed order n and bandwidths whose relevant entries are randomly generated integer values from a prescribed interval. The band matrix B is TT^T , and the program

calculates the maximal eigenvalues of $Au = \lambda Bu$. For the sequence of values of $p, p = 2, \dots, \min([n/5], 10)$, TESTB exercises SIMITZ for successive values of $em, em = 1, \dots, p - 1$. For each value of $i, i = 1, \dots, em$, TESTB computes the residuals $Ax_i - d_i Bx_i$ and their Euclidean norms relative to the standard inner product. Each norm is normalized by the difference $|d_i| - e$, and for each value of em the quantity $\max_{1 \leq i \leq em} \|Ax_i - d_i Bx_i\| / (|d_i| - e)$, the value k of i for which the maximum occurs, and the corresponding geometric mean with unit weights are listed. Also listed are the relevant nonzero diagonals of A and T and the final eigenvalues computed for $em = p - 1$.

Figure 1 shows an output listing from the executable program TESTB on the CDC 6600 under the NOS/BE operating system and Fortran Extended, Version 4. Here A and B are of order 30 and each of bandwidth 7 having relevant entries between -99 and $+99$. Listed are the main diagonal and the three adjacent lower diagonals of T and A beginning with the entries in the first column. Here $eps = 10^{-10}$ and $km = 100$. Note how the relative nearness in magnitude of the first three eigenvalues inhibits the convergence of the second eigenvector when $p = 3$ and $em = 2$ resulting in acceptance of the first eigenvector only. The fourth eigenvalue, however, is in absolute value far enough away from this cluster to permit successful convergence when $p = 4$ and $em = 1, 2$, and 3 . This phenomenon points to a procedure for pursuing a solution when p is initially chosen too small. SIMITZ may be reentered with X containing the approximate eigenvectors calculated for the smaller value of p as initial approximations for use with p increased in size. Significant processor time may often be saved in this fashion.

ADJACENT NON-ZERO LOWER DIAGONALS OF T, TT' = B														
53	86	78	60	85	99	73	64	80	88	93	93	57	52	83
74	99	88	87	52	64	86	95	66	86	55	95	61	76	70
-26	22	49	65	-30	38	82	-9	-87	56	18	95	93	-46	-31
1	90	-44	75	-6	-49	99	90	-62	16	2	-40	-67	14	
-25	6	54	32	-73	-35	-1	-46	6	-13	-77	77	-51	-85	31
-24	88	-66	-18	-41	5	80	74	-80	-2	-6	78	-97		
-76	19	-26	-5	97	70	6	-95	45	-30	-12	-65	-92	9	39
97	-19	26	86	-48	-52	-31	49	-8	-85	5	20			
ADJACENT NON-ZERO LOWER DIAGONALS OF A = A'														
-3	50	-55	-73	-47	88	-33	-28	-77	-40	96	78	9	-54	-39
10	95	-53	-80	0	52	44	0	-50	-94	-22	29	-36	-4	36
0	65	59	-29	-84	-4	41	-90	72	-65	74	36	82	-86	-5
22	56	94	9	-47	84	-97	-1	57	-22	-1	-55	-15	51	
62	5	-41	-95	-47	-91	14	84	-83	69	43	-27	61	-79	14
21	85	-5	-61	-16	-15	90	-44	-22	-66	50	-34	-97		
-17	-94	78	-95	-89	-31	-43	52	-64	-70	71	95	-69	-12	87
91	-93	6	39	66	32	-20	92	-11	0	-98	-12			
P	EM (IN)	EM (OUT)	KS	K	MAX RESIDUAL	MEAN RESIDUAL	TIME (SECS)							
2	1	1	30	1	3.70E-11	3.70E-11	.42							
3	1	1	26	1	2.81E-11	2.81E-11	.62							
3	2	1	56	1	3.00E-11	3.00E-11	1.00							
4	1	1	15	1	2.17E-11	2.17E-11	.69							
4	2	2	15	2	2.01E-11	1.81E-11	.69							
4	3	3	24	3	2.53E-11	2.09E-11	.92							
5	1	1	13	1	2.20E-11	2.20E-11	1.02							
5	2	2	13	1	2.03E-11	1.99E-11	1.03							
5	3	3	13	1	2.24E-11	1.60E-11	1.01							
5	4	4	33	1	2.37E-11	6.67E-12	1.44							
6	1	1	13	1	2.09E-11	2.09E-11	1.36							
6	2	2	13	2	2.16E-11	2.04E-11	1.36							
6	3	3	13	2	2.89E-11	2.03E-11	1.33							
6	4	4	21	4	2.02E-10	3.63E-11	1.76							
6	5	4	31	1	2.64E-11	1.09E-11	2.10							
FINAL EIGENVALUES														
2.393019487013E+00	-1.245936323012E+00	1.188248083996E+00	-2.273759446561E-01	-9.705562962258E-02										

Fig. 1

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ALGORITHM

```

SUBROUTINE SIMITZ(N, P, KM, EPS, IP, OP, INF, EM, X, MN, D, WK)      10
C*****                                                              20
C IDENTIFICATION                                                    30
C   SIMITZ - ITERATIVE COMPUTATION OF EIGENVALUES LARGEST IN MAGNI-  40
C   TUDE AND CORRESPONDING EIGENVECTORS OF A REAL GENERAL-      50
C   IZED SYMMETRIC MATRIX                                          60
C   FORTRAN SUBROUTINE SUBPROGRAM                                  70
C   US AIR FORCE FLIGHT DYNAMICS LABORATORY                        80
C   WRIGHT-PATTERSON AFB, OHIO 45433                              90
C PURPOSE                                                           100
C   A REAL N-SQUARE MATRIX C IS B-SYMMETRIC RELATIVE TO AN N-SQUARE  110
C   POSITIVE DEFINITE MATRIX B IN CASE BC = (C-TRANPOSED)B.      120
C   GIVEN AS OPTIONAL INPUT A SET OF P INITIAL APPROXIMATE      130
C   EIGENVECTORS OF A REAL N-SQUARE B-SYMMETRIC MATRIX C CORRES-  140
C   PONDING TO P EIGENVALUES OF C LARGEST IN MAGNITUDE, SIMITZ COM-  150
C   PUTES EM EIGENVALUES AND EM CORRESPONDING EIGENVECTORS TO A   160
C   PRECISION DEPENDENT ON THE STRUCTURE OF B AND C AND ON A GIVEN  170
C   TOLERANCE EPS. THE MATRIX B IS PRESENTED TO SIMITZ AS AN ALGO-  180
C   RITHM FOR CALCULATING THE STANDARD INNER PRODUCT (W, BZ) =    190
C   (W-TRANPOSED)BZ GIVEN COLUMN N-VECTORS W AND Z IMPLEMENTED AS  200
C   A FORTRAN COMPATIBLE REAL FUNCTION SUBPROGRAM. THE MATRIX C IS  210
C   PRESENTED AS A SUBROUTINE SUBPROGRAM WHICH GIVEN A COLUMN     220
C   N-VECTOR Z CALCULATES ITS IMAGE W = CZ UNDER THE MATRIX C.    230
C   DEPENDING ON THE CHOICE OF B AND C, SIMITZ APPLIES TO A WIDE   240
C   VARIETY OF SYMMETRIC EIGENPROBLEMS.                            250
C CONTROL                                                           260
C   DIMENSION X(MN,P), D(P), WK(K)                                270
C   INTEGER P, EM                                                  280
C   REAL IP                                                         290
C   EXTERNAL IP, INF, OP                                           300
C   .                                                                310
C   .                                                                320
C   .                                                                330
C   .                                                                340
C   CALL SIMITZ(N, P, KM, EPS, IP, OP, INF, EM, X, MN, D, WK)     350
C   WHERE                                                           360
C   N IS AN INTEGER INPUT VARIABLE, THE ORDER OF THE MATRIX C.    370
C   P IS AN INTEGER INPUT VARIABLE, THE NUMBER OF SIMULTANEOUS     380
C   ITERATION VECTORS.                                             390
C   KM AS AN INTEGER INPUT VARIABLE IS IN MAGNITUDE THE MAXIMUM    400
C   NUMBER OF ITERATION STEPS TO BE EXECUTED. IF KM IDENTIFIES    410
C   A NEGATIVE VALUE THEN P INITIAL APPROXIMATE EIGENVECTORS     420
C   ARE ASSUMED TO BE PRESENT IN THE ARRAY X. OTHERWISE SIMITZ    430
C   SUPPLIES RANDOM INITIAL EIGENVECTORS.                          440
C   KM AS AN INTEGER OUTPUT VARIABLE IDENTIFIES THE NUMBER KS OF   450
C   ITERATION STEPS FINALLY USED IN THE CALCULATION OF EM         460
C   EIGENVECTORS.                                                 470
C   EPS IS A REAL INPUT VARIABLE, THE TOLERANCE FOR ACCEPTING     480
C   EIGENVECTORS. AS SOON AS SUCCESSIVE ITERATES OF THE RITZ     490
C   500

```

C VALUES ABS(D(H+1)) DIFFER BY LESS THAN ABS(D(H+1))*EPS/10.0 510
 C THEN D(H+1) IS ACCEPTED AS AN EIGENVALUE AND H, THE NUMBER 520
 C OF PREVIOUSLY ACCEPTED EIGENVALUES, IS INCREASED BY 1. AS 530
 C SOON AS THE ERROR QUANTITIES F(I), NORMALIZED RESIDUALS, 540
 C SATISFY $D(I)*F(I)/(D(I) - D(P))$.LT. EPS, THEN G, THE NUM- 550
 C BER OF ALREADY ACCEPTED RITZ VECTORS, IS INCREASED TO 560
 C $G + 1$, $I = G + 1, \dots, L$. THE F(I) ARE DISCOUNTED WITH 570
 C SUCCESSIVE ITERATIONS TO FORCE CONVERGENCE IN CASE OF UN- 580
 C FORTUNATE CHOICE OF PARAMETERS. IF M SIGNIFICANT DIGITS 590
 C OF ACCURACY ARE REQUIRED OF THE EIGENVALUES, THEN SET 600
 C EPS EQUAL TO $10.0^{*(-M)}$ AS A GENERAL RULE. 610
 C IP IS AN EXTERNAL INPUT VARIABLE, THE NAME OF A FORTRAN COM- 620
 C PATIBLE REAL FUNCTION SUBPROGRAM OF THE FORM IP(N, Z, W) 630
 C WHICH MUST RETURN THE INNER PRODUCT $(W, BZ) =$ 640
 C $(W-TRANPOSED)BZ$ OF THE VECTORS IDENTIFIED BY THE N-ARRAYS 650
 C Z AND W RELATIVE TO THE POSITIVE DEFINITE MATRIX B. 660
 C OP IS AN EXTERNAL INPUT VARIABLE, THE NAME OF A FORTRAN COM- 670
 C PATIBLE SUBROUTINE SUBPROGRAM OF THE FORM OP(N, Z, W) 680
 C WHICH MUST CALCULATE THE IMAGE W OF THE VECTOR IDENTIFIED 690
 C BY THE N-ARRAY Z UNDER THE N-SQUARE MATRIX C WITHOUT OVER- 700
 C WRITING Z. 710
 C INF IS AN EXTERNAL INPUT VARIABLE, THE NAME OF A FORTRAN COM- 720
 C PATIBLE SUBROUTINE SUBPROGRAM WHICH MAY BE USED FOR 730
 C OBTAINING INFORMATION OR TO EXERT CONTROL DURING EXECUTION 740
 C OF SIMITZ. INF HAS THE FORM INF(KS, G, H, F) WHERE 750
 C KS IS AN INTEGER OUTPUT VARIABLE, THE NUMBER OF THE NEXT 760
 C ITERATION STEP. 770
 C G IS AN INTEGER OUTPUT VARIABLE, THE NUMBER OF ALREADY 780
 C ACCEPTED EIGENVECTORS. 790
 C H IS AN INTEGER OUTPUT VARIABLE, THE NUMBER OF ALREADY 800
 C ACCEPTED EIGENVALUES. 810
 C F IS A REAL OUTPUT VARIABLE P-ARRAY, ERROR QUANTITIES 820
 C MEASURING RESPECTIVELY THE STATE OF CONVERGENCE OF 830
 C THE P SIMULTANEOUS ITERATION VECTORS. IN ADDITION, 840
 C IF CONVERGENCE FAILS IN SUBROUTINE IMTQL2 AFTER G 850
 C EIGENVECTORS HAVE BEEN ACCEPTED, THEN F(G+1) IS RE- 860
 C PLACED BY $1000.*FLOAT(IERR)$ WHERE IERR IS THE ERROR 870
 C INDICATOR OUTPUT BY IMTQL2. EACH ELEMENT OF THE ARRAY 880
 C F IS INITIALLY SET BY SIMITZ TO THE VALUE 4.0. 890
 C EM AS AN INTEGER INPUT VARIABLE IS THE NUMBER OF EIGENVALUES 900
 C TO BE COMPUTED, 0 .LT. EM .LT. P .LE. N .LE. MN. 910
 C EM AS AN INTEGER OUTPUT VARIABLE IS THE NUMBER OF EIGENVECTORS 920
 C COMPUTED THROUGH KM ITERATION STEPS. 930
 C X AS A REAL N-BY-P INPUT ARRAY IS A SET OF P OPTIONAL INITIAL 940
 C APPROXIMATE EIGENVECTORS $X(I,1), \dots, X(I,P)$, $I = 1, \dots,$ 950
 C N, INTERPRETED BY SIMITZ IF KM IS NEGATIVE. 960
 C X AS A REAL N-BY-P OUTPUT ARRAY IS A SET OF EM EIGENVECTORS 970
 C $X(I,1), \dots, X(I,EM)$, $I = 1, \dots, N$, COMPUTED THROUGH 980
 C IABS(KM) ITERATION STEPS WITH THE REMAINDER OF X CONSISTING 990
 C OF P - EM APPROXIMATE EIGENVECTORS. THE N-BY-P MATRIX X 1000
 C SATISFIES $(X-TRANPOSED)BX = I$, THAT IS, THE EIGENVECTORS 1010
 C OF C ARE B-ORTHONORMAL. 1020
 C MN IS AN INTEGER INPUT VARIABLE WHICH IDENTIFIES THE LEADING 1030
 C DIMENSION IN THE CALLING PROGRAM OF THE ARRAY X. 1040
 C D IS A REAL OUTPUT P-ARRAY OF WHICH $D(1), \dots, D(EM)$ ARE THE 1050
 C EIGENVALUES OF C LARGEST IN MAGNITUDE IN DECREASING ORDER 1060
 C CORRESPONDING TO THE COMPUTED EIGENVECTORS $X(I,1), \dots,$ 1070
 C $X(I,EM)$, $I = 1, \dots, N$. $D(EM+1), \dots, D(P-1)$ CONTAIN 1080
 C APPROXIMATIONS TO PROGRESSIVELY SMALLER SUCH EIGENVALUES. 1090
 C $D(P)$ CONTAINS THE MOST RECENTLY COMPUTED VALUE OF E, WHERE 1100
 C THE INTERVAL $(-E, E)$ IS THE INTERVAL OVER WHICH THE 1110
 C CHEBYSHEV ACCELERATION WAS PERFORMED. 1120
 C WK THE INITIAL LOCATION OF AT LEAST $P**2 + 3*P + 2*N = K$ 1130
 C CONSECUTIVE STORAGE LOCATIONS WHICH MAY NOT BE OVER- 1140
 C WRITTEN WHILE SIMITZ IS IN EXECUTION. 1150
 C OTHER PROGRAMMING INFORMATION 1160
 C SIMITZ EMPLOYS A DATA STATEMENT TO ASSIGN TO A MACHINE DEPEND- 1170
 C ENT REAL VARIABLE MT THE QUOTIENT OF THE SMALLEST POSITIVE 1180
 C REAL VALUE REPRESENTABLE BY FORTRAN AND THE SMALLEST REAL VALUE 1190
 C WHOSE SUM WITH 1.0 EXCEEDS 1.0. 1200
 C 1210
 C THE PERFORMANCE OF SIMITZ IS STRONGLY DEPENDENT UPON THE CHOICE 1220
 C OF INPUT PARAMETERS AND UPON THE CAREFUL PREPARATION OF THE 1230
 C SUBPROGRAMS IP AND OP. THE USER SHOULD CONSIDER USING HIS OWN 1240
 C ACTIVE SUBROUTINE INF TO MONITOR PROGRESS OF SIMITZ RELATIVE TO 1250
 C HIS CHOICE OF INPUT PARAMETERS IF NO INFORMATION IS OTHERWISE 1260
 C AVAILABLE CONCERNING THE LOCATIONS OF THE RELEVANT EIGENVALUES. 1270
 C RECALL THAT SIMITZ MAY BE REENTERED WITH KM .LT. 0 WITHOUT LOSS 1280


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C      OF INFORMATION TO PERMIT CONSERVATIVE INITIAL CHOICES OF      1290
C      ABS(KM), EPS AND P.                                          1300
C      OTHER PROGRAMS REQUIRED                                       1310
C      FUNCTION RANF                                               1320
C      RETURNS UNIFORMLY DISTRIBUTED RANDOM NUMBERS ON THE OPEN    1330
C      INTERVAL (0, 1) GIVEN ANY ONE ARGUMENT OF ANY TYPE.        1340
C      SUBROUTINE TRED2                                             1350
C      IS THE EISPACK (4) PROGRAM WHICH COMPUTES A HOUSEHOLDER     1360
C      TRIDIAGONAL FORM OF A REAL SYMMETRIC MATRIX.               1370
C      SUBROUTINE IMTQL2                                            1380
C      IS THE EISPACK PROGRAM WHICH COMPUTES THE EIGENVALUES AND   1390
C      ORTHONORMAL EIGENVECTORS OF A SYMMETRIC TRIDIAGONAL MATRIX. 1400
C      FUNCTION IP                                                 1410
C      IS DESCRIBED ABOVE.                                         1420
C      SUBROUTINE OP                                               1430
C      IS DESCRIBED ABOVE.                                         1440
C      SUBROUTINE INF                                              1450
C      IS DESCRIBED ABOVE.                                         1460
C      METHOD                                                       1470
C      SIMITZ REPRESENTS RESULTS OF EXTENSIVE MODIFICATIONS AND TESTS 1480
C      OF SUBROUTINE RITZIT (1), AN ANSI FORTRAN TRANSLATION OF THE 1490
C      ALGOL 60 PROCEDURE OF THE SAME NAME (3). THE BASIC RUTISHAUSER 1500
C      -REINSCH ALGORITHM IS PRESERVED.                            1510
C      REFERENCES                                                  1520
C      (1) PAUL J. NIKOLAI AND NAI-KUAN TSAO, THE ARL LINEAR ALGEBRA 1530
C      LIBRARY HANDBOOK, ARL TR 74-0106, AEROSPACE RESEARCH LABOR- 1540
C      ATORIES, WRIGHT-PATTERSON AFB, OHIO, 1974.                 1550
C      (2) HEINZ RUTISHAUSER, COMPUTATIONAL ASPECTS OF F.L. BAUER S 1560
C      SIMULTANEOUS ITERATION METHOD, NUMER. MATH. 13(1969), 4-13. 1570
C      (3) -----, SIMULTANEOUS ITERATION METHOD FOR SYM-        1580
C      METRIC MATRICES, NUMER. MATH. 16(1970), 205-223.           1590
C      (4) B.T. SMITH ET AL, MATRIX EIGENSYSTEM ROUTINES-EISPACK   1600
C      GUIDE, LECTURE NOTES IN COMPUTER SCIENCE 6, SPRINGER-VERLAG 1610
C      NEW YORK, 1974.                                           1620
C      PLEASE REFER QUESTIONS, COMMENTS OR SUGGESTIONS TO        1630
C      PAUL J. NIKOLAI                                           1640
C      AFFDL/FBR                                                 1650
C      WRIGHT-PATTERSON AFB, OH 45433                             1660
C      (513)-255-5350                                           1670
C      *****                                                    1680
C      EXTERNAL   INF,   IP,   OP                                 1690
C      INTEGER   EM,   G,   H,   I,   IG,   IK,   J,             1700
C      *         JK,   JP,   K,   KM,   KS,   L,   LF,           1710
C      *         LI,   M,   MN,   MI,   N,   P,   ZI,            1720
C      *         Z2                                             1730
C      LOGICAL   ORIG                                           1740
C      REAL     D,   E,   EPS,   E1,   E2,   IP,   MT,          1750
C      *         S,   T,   WK,   X                               1760
C      DIMENSION X(MN,1), D(1), WK(P,P,1)                        1770
C      DATA MT /.220360641585062E-279/                          1780
C      THE LOCAL VARIABLE ARRAYS FROM (3) ARE ASSIGNED TO THE    1790
C      VARIABLE ARRAY WK AS FOLLOWS                              1800
C      WK(I,J,1) = B(I,J), I, J = 1, ..., P.                    1810
C      WK(I,1,2) = CX(I), I = 1, ..., P.                         1820
C      WK(I,2,2) = F(I), I = 1, ..., P.                          1830
C      WK(I,3,2) = RQ(I), I = 1, ..., P.                         1840
C      WK(I,4,2) = U(I), I = 1, ..., N.                          1850
C      WK(I+N,4,2) = W(I), I = 1, ..., N.                       1860
C      NOT NEEDED ARE V(I), I = 1, ..., N, R(I), I = 1, ..., P, AND 1870
C      Q(I,J), I, J = 1, ..., P.                                1880
C      THE NEXT STATEMENT IS START.                              1890
C      E = .0E+00                                               1900
C      G = 0                                                    1910
C      IG = 1                                                  1920
C      H = 0                                                    1930
C      Z1 = 0                                                  1940
C      Z2 = 0                                                  1950
C      KS = 0                                                  1960
C      M = 1                                                  1970
C      WK(P,1,2) = .0E+00                                       1980

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      DO 10 L = 1, P                                2060
        WK(L,2,2) = .4E+01                          2070
        WK(L,3,2) = .0E+00                          2080
10    CONTINUE                                    2090
      IF (KM) 50, 50, 20                            2100
20    DO 40 L = 1, P                                2110
      DO 30 J = 1, N                                2120
        X(J,L) = .2E+01*RANF(.2E+01) - .1E+01      2130
30    CONTINUE                                    2140
40    CONTINUE                                    2150
50    KM = IABS(KM)                                2160
      ASSIGN 60 TO IK                              2170
      LF = IG                                       2180
      LI = P                                        2190
      GO TO 990                                     2200
C      RAYLEIGH-RITZ STEP                          2210
C      STATEMENT 60 IS LOOP.                       2220
60    DO 80 K = IG, P                              2230
      CALL OP(N, X(1,K), WK(1,4,2))                2240
      DO 70 J = 1, N                                2250
        X(J,K) = WK(J,4,2)                         2260
70    CONTINUE                                    2270
80    CONTINUE                                    2280
      ASSIGN 90 TO IK                              2290
      LF = IG                                       2300
      LI = P                                        2310
      GO TO 990                                     2320
90    IF (KS) 150, 100, 150                        2330
C      MEASURES AGAINST UNHAPPY CHOICE OF INITIAL VECTORS 2340
100   DO 130 K = 1, P                              2350
      IF (WK(K,K,1)) 130, 110, 130                 2360
110   DO 120 I = 1, N                                2370
      X(I,K) = .2E+01*RANF(.2E+01) - .1E+01      2380
120   CONTINUE                                    2390
      KS = 1                                        2400
130   CONTINUE                                    2410
      IF (KS - 1) 150, 140, 150                    2420
140   ASSIGN 60 TO IK                              2430
      LF = 1                                        2440
      LI = P                                        2450
      GO TO 990                                     2460
150   DO 180 K = IG, P                              2470
      DO 170 L = K, P                                2480
        S = .0E+00                                  2490
        DO 160 I = L, P                              2500
          S = S + WK(I,K,1)*WK(I,L,1)              2510
160   CONTINUE                                    2520
      WK(L,K,1) = -S                                2530
170   CONTINUE                                    2540
180   CONTINUE                                    2550
      CALL TRED2(P, P - G, WK(IG,IG,1), D(IG), WK(1,4,2), WK(IG,IG,1)) 2560
      CALL IMTQL2(P, P - G, D(IG), WK(1,4,2), WK(IG,IG,1), L) 2570
      WK(IG,2,2) = AMAX1(WK(IG,2,2), .1E+04*FLOAT(L)) 2580
      DO 190 K = IG, P                              2590
        D(K) = SQRT(AMAX1(-D(K), .0E+00))          2600
190   CONTINUE                                    2610
C      REORDERING EIGENVALUES AND EIGENVECTORS ACCORDING TO SIZE OF 2620
C      THE FORMER IS ACCOMPLISHED IN SUBROUTINE IMTQL2. 2630
      DO 230 J = 1, N                                2640
      DO 210 K = IG, P                              2650
        S = .0E+00                                  2660
        DO 200 L = IG, P                              2670
          S = S + X(J,L)*WK(L,K,1)                 2680
200   CONTINUE                                    2690
      WK(K,4,2) = S                                 2700
210   CONTINUE                                    2710
      DO 220 K = IG, P                              2720
        X(J,K) = WK(K,4,2)                         2730
220   CONTINUE                                    2740
230   CONTINUE                                    2750
      KS = KS + 1                                  2760
      E = AMAX1(D(P), E)                            2770
C      RANDOMIZATION                               2780
      IF (3 - Z1) 260, 240, 240                    2790
240   DO 250 J = 1, N                                2800
      X(J,P) = .2E+01*RANF(.2E+01) - .1E+01      2810

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250	CONTINUE	2820
	JP = P - 1	2830
	ASSIGN 260 TO IK	2840
	LF = P	2850
	L1 = P	2860
	GO TO 990	2870
C	COMPUTE CONTROL QUANTITIES CX(I).	2880
260	DO 310 K = IG, JP	2890
	S = (D(K) - E)*(D(K) + E)	2900
	IF (S) 270, 270, 280	2910
270	WK(K,1,2) = .0E+00	2920
	GO TO 310	2930
280	IF (E) 300, 290, 300	2940
290	WK(K,1,2) = .1E+04 + ALOG(D(K))	2950
	GO TO 310	2960
300	WK(K,1,2) = ALOG((D(K) + SQRT(S))/E)	2970
310	CONTINUE	2980
C	ACCEPTANCE TEST FOR EIGENVALUES INCLUDING ADJUSTMENT OF EM AND	2990
C	H SUCH THAT D(EM) .GT. E, D(H) .GT. E AND D(EM) DOES NOT	3000
C	OSCILLATE STRONGLY	3010
	I = Z1 - 1	3020
	K = G	3030
320	K = K + 1	3040
	IF (EM - K) 370, 330, 330	3050
330	IF (D(K) - E) 360, 360, 340	3060
340	IF (I) 320, 320, 350	3070
350	IF (D(K) - .999E+00*WK(K,3,2)) 360, 360, 320	3080
360	CONTINUE	3090
	EM = K - 1	3100
C	STATEMENT 370 IS EX4.	3110
370	IF (EM) 380, 1130, 380	3120
380	K = H	3130
	S = .1E+01 + .1E+00*EPS	3140
390	K = K + 1	3150
	IF (D(K)) 400, 410, 400	3160
400	IF (D(K) - S*WK(K,3,2)) 390, 390, 410	3170
410	CONTINUE	3180
	H = K - 1	3190
	K = EM	3200
420	K = K + 1	3210
	IF (K - H) 430, 430, 450	3220
430	IF (D(K) - E) 440, 440, 420	3230
440	CONTINUE	3240
	H = K - 1	3250
C	ACCEPTANCE TEST FOR EIGENVECTORS	3260
450	L = G	3270
	E2 = .0E+00	3280
	DO 590 K = IG, JP	3290
	IF (K - (L + 1)) 510, 460, 510	3300
C	CHECK FOR NESTED EIGENVALUES	3310
460	L = K	3320
	L1 = K	3330
	S = .5E+00/FLOAT(KS)	3340
	T = .1E+01/FLOAT(KS*M)	3350
470	L = L + 1	3360
	IF (L - JP) 480, 480, 490	3370
480	IF (WK(L,1,2)*(WK(L,1,2) + S) + T - WK(L-1,1,2)*WK(L-1,1,2))	3380
*	490, 490, 470	3390
490	CONTINUE	3400
	L = L - 1	3410
C	THE NEXT STATEMENT IS EX5.	3420
	IF (L - H) 510, 510, 500	3430
500	L = L1 - 1	3440
	GO TO 600	3450
510	CALL OP(N, X(1,K), WK(1,4,2))	3460
	S = .0E+00	3470
	DO 540 J = 1, L	3480
	IF (ABS(D(J) - D(K)) - .1E-01*D(K)) 520, 540, 540	3490
520	T = IP(N, WK(1,4,2), X(1,J))	3500
	DO 530 I = 1, N	3510
	WK(I,4,2) = WK(I,4,2) - T*X(I,J)	3520
530	CONTINUE	3530
	S = S + T*T	3540
540	CONTINUE	3550
	T = .1E+01	3560
	IF (S .NE. .0E+00) T = IP(N, WK(1,4,2), WK(1,4,2))	3570

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      E2 = AMAX1(E2, SQRT(T/(S + T)))
      IF (K - L) 590, 550, 590
C     TEST FOR ACCEPTANCE OF GROUP OF EIGENVECTORS
550   IF (L .GE. EM .AND. D(EM)*WK(EM,2,2) .LT. EPS*(D(EM) - E))
      *   G = EM
      IF (E2 - WK(L,2,2)) 560, 580, 580
560   DO 570 J = L1, L
      WK(J,2,2) = E2
570   CONTINUE
580   IF (L .LE. EM .AND. D(L)*WK(L,2,2) .LT. EPS*(D(L) - E)) G = L
590   CONTINUE
C     ADJUST M.
C     STATEMENT 600 IS EX6.
600   IG = G + 1
      IF (E - .4E-01*D(1)) 610, 610, 620
610   M = 1
      K = 1
      GO TO 630
620   E2 = .2E+01/E
      E1 = .51E+00*E2
      K = 2*INT(.4E+01/AMIN1(WK(1,1,2), .4E+01))
      M = MIN0(M, K)
C     REDUCE EM IF CONVERGENCE WOULD BE TOO SLOW.
630   IF (WK(EM,2,2)) 640, 690, 640
640   IF (FLOAT(KS) - .9E+00*FLOAT(KM)) 650, 690, 690
650   S = FLOAT(K)*WK(EM,1,2)
      IF (S - .5E-01) 660, 670, 670
660   T = .5E+00*S*WK(EM,1,2)
      GO TO 680
670   T = WK(EM,1,2) + ALOG(.5E+00 + .5E+00*EXP(-.2E+01*S))/FLOAT(K)
680   S = ALOG(D(EM)*WK(EM,2,2)/(EPS*(D(EM) - E)))
      IF (S*FLOAT(KS) .GT. T*FLOAT((KM - KS)*KM)) EM = EM - 1
C     STATEMENT 690 IS EX2.
690   DO 700 K = IG, JP
      WK(K,3,2) = D(K)
700   CONTINUE
      CALL INF(KS, G, H, WK(1,2,2))
      IF (G .GE. EM .OR. KS .GE. KM) GO TO 1130
C     STATEMENT 710 IS EX1.
710   IF (KS + M - KM) 730, 730, 720
720   Z2 = -1
      IF (M .GT. 1) M = 2*((KM - KS + 1)/2)
730   M1 = M
C     SHORTCUT LAST INTERMEDIATE BLOCK IF ALL F(I) ARE SUFFICIENTLY
C     SMALL.
      IF (L - EM) 780, 740, 740
740   S = D(EM)*WK(EM,2,2)/(EPS*(D(EM) - E))
      T = S*S - .1E+01
      IF (T) 60, 60, 750
750   S = ALOG(S + SQRT(T))/(WK(EM,1,2) - WK(H+1,1,2))
      M1 = 2*INT(.5E+00*S + .101E+01)
      IF (M1 - M) 770, 770, 760
760   M1 = M
      GO TO 780
770   Z2 = -1
C     CHEBYSHEV ITERATION
780   IF (M - 1) 900, 790, 820
790   DO 810 K = IG, P
      CALL OP(N, X(1,K), WK(1,4,2))
      DO 800 I = 1, N
      X(I,K) = WK(I,4,2)
800   CONTINUE
810   CONTINUE
      GO TO 900
820   L1 = M1 - 4
      DO 890 K = IG, P
      CALL OP(N, X(1,K), WK(1,4,2))
      DO 830 I = 1, N
      IK = I + N
      WK(IK,4,2) = E1*WK(I,4,2)
830   CONTINUE
      CALL OP(N, WK(N+1,4,2), WK(1,4,2))
      DO 840 I = 1, N
      X(I,K) = E2*WK(I,4,2) - X(I,K)
840   CONTINUE
      IF (L1) 890, 850, 850
850   DO 880 J = 4, M1, 2

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      CALL OP(N, X(1,K), WK(1,4,2))
      DO 860 I = 1, N
        IK = I + N
        WK(IK,4,2) = E2*WK(I,4,2) - WK(IK,4,2)
860    CONTINUE
      CALL OP(N, WK(N+1,4,2), WK(1,4,2))
      DO 870 I = 1, N
        X(I,K) = E2*WK(I,4,2) - X(I,K)
870    CONTINUE
880    CONTINUE
890    CONTINUE
900    ASSIGN 910 TO IK
      LF = IG
      LI = P
      GO TO 990
C      DISCOUNTING THE ERROR QUANTITIES F
910    IF (G - H) 920, 970, 970
920    IF (M - 1) 950, 930, 950
930    DO 940 K = IG, H
      WK(K,2,2) = WK(K,2,2)*(D(H+1)/D(K))
940    CONTINUE
      GO TO 970
950    T = EXP(-FLOAT(M1)*WK(H+1,1,2))
      DO 960 K = IG, H
        S = EXP(-FLOAT(M1)*(WK(K,1,2) - WK(H+1,1,2)))
        WK(K,2,2) = S*WK(K,2,2)*(.1E+01 + T*T)/(.1E+01 + (S*T)*(S*T))
960    CONTINUE
970    KS = KS + M1
      Z2 = Z2 - M1
C      POSSIBLE REPETITION OF INTERMEDIATE STEPS
      IF (Z2) 980, 710, 710
980    Z1 = Z1 + 1
      Z2 = 2*Z1
      M = M + M
      GO TO 60
C      PERFORMS ORTHONORMALIZATION OF COLUMNS 1 THROUGH LI OF ARRAY
C      X ASSUMING THAT COLUMNS 1 THROUGH LF - 1 ARE ALREADY ORTHO-
C      NORMAL
990    DO 1120 K = LF, LI
      ORIG = .TRUE.
1000    T = .0E+00
      JK = K - 1
      IF (JK) 1040, 1040, 1010
1010    DO 1030 I = 1, JK
      S = IP(N, X(1,I), X(1,K))
      IF (ORIG) WK(K,I,1) = S
      T = T + S*S
      DO 1020 J = 1, N
        X(J,K) = X(J,K) - S*X(J,I)
1020    CONTINUE
1030    CONTINUE
1040    S = IP(N, X(1,K), X(1,K))
      T = S + T
      IF (S - .1E-01*T) 1060, 1060, 1050
1050    IF (T - MT) 1060, 1060, 1080
1060    ORIG = .FALSE.
      IF (S - MT) 1070, 1070, 1000
1070    S = .0E+00
1080    S = SQRT(S)
      WK(K,K,1) = S
      IF (S) 1090, 1100, 1090
1090    S = .1E+01/S
1100    DO 1110 J = 1, N
      X(J,K) = S*X(J,K)
1110    CONTINUE
1120    CONTINUE
      GO TO IK, (60, 90, 260, 910, 1140)
C      STATEMENT 1130 IS EX.
1130    EM = G
C      SOLVE EIGENVALUE PROBLEM OF PROJECTION OF MATRIX C.
      ASSIGN 1140 TO IK
      LF = 1
      LI = JP
      GO TO 990
1140    DO 1160 K = 1, JP
      CALL OP(N, X(1,K), X(1,P))

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DO 1150 I = 1, K
WK(K,I,1) = -IP(N, X(1,I), X(1,P))
1150 CONTINUE
1160 CONTINUE
CALL TRED2(P, JP, WK, D, WK(1,4,2), WK)
CALL IMTQL2(P, JP, D, WK(1,4,2), WK, L)
WK(IG,2,2) = AMAX1(WK(IG,2,2), .1E+04*FLOAT(L))
C ARRANGE EIGENVALUES IN ORDER OF DECREASING ABSOLUTE VALUE.
DO 1210 J = 1, JP
K = J
DO 1170 I = J, JP
IF (ABS(D(I)) .GT. ABS(D(K))) K = I
1170 CONTINUE
IF (K - J) 1200, 1200, 1180
1180 T = D(K)
D(K) = D(J)
D(J) = T
DO 1190 I = 1, JP
T = WK(I,K,1)
WK(I,K,1) = WK(I,J,1)
WK(I,J,1) = T
1190 CONTINUE
1200 D(J) = -D(J)
1210 CONTINUE
DO 1250 J = 1, N
DO 1230 I = 1, JP
S = .0E+00
DO 1220 K = 1, JP
S = S + X(J,K)*WK(K,I,1)
1220 CONTINUE
WK(I,4,2) = S
1230 CONTINUE
DO 1240 I = 1, JP
X(J,I) = WK(I,4,2)
1240 CONTINUE
1250 CONTINUE
KM = KS
D(P) = E
RETURN
END
C PROGRAM TESTB(INPUT, OUTPUT, TAPE5 = INPUT, TAPE6 = OUTPUT)
DIMENSION A(50,6), T(50,6), Y(50)
COMMON NT, T, NA, A, ND, Y
INTEGER Y
DIMENSION X(50,10), D(10), WK(115,2)
EXTERNAL INTROS, ABAND, BAND
C
C THE VARIABLE ND MUST IDENTIFY A VALUE AT LEAST EQUAL TO N.
C THE VALUE ASSIGNED TO THE VARIABLE ND MUST AGREE WITH THE LEADING
C DIMENSION OF THE VARIABLES A, T, X, AND Y. THE LEADING DIMENSION OF
C THE VARIABLE WK MUST EQUAL OR EXCEED ND + 65. THE SECOND DIMENSION
C OF A AND T MUST EXCEED NA AND NT, RESPECTIVELY. OTHERWISE RECOMPILE.
C NOTE THAT DIMENSIONS OF THE VARIABLES A, T, AND Y MUST AGREE WITH
C THEIR COUNTERPARTS IN FUNCTION BAND AND SUBROUTINE ABAND.
C
C N IS AN INTEGER INPUT VARIABLE, THE ORDER OF THE MATRICES A AND
C B. IF N .LE. 0, PROCESSING CEASES.
C NA IS AN INTEGER INPUT VARIABLE, THE NUMBER OF ADJACENT NON-ZERO
C DIAGONALS STRICTLY BELOW THE MAIN DIAGONAL OF THE MATRIX A.
C NT IS AN INTEGER INPUT VARIABLE, THE NUMBER OF ADJACENT NON-ZERO
C DIAGONALS STRICTLY BELOW THE MAIN DIAGONAL OF THE MATRIX T.
C KM IS AN INTEGER INPUT VARIABLE, THE MAXIMUM NUMBER OF ITERATION
C STEPS TO BE EXECUTED BY SIMITZ ON EACH CALL. IF KM IDENTIFIES
C A NEGATIVE VALUE, SIMITZ WILL USE PREVIOUSLY COMPUTED EIGENVEC
C TORS AS INITIAL APPROXIMATIONS FOR A GIVEN VALUE OF P WHILE EM
C IS INCREASED.
C EPS IS A REAL INPUT VARIABLE, THE TOLERANCE FOR ACCEPTING EIGEN-
C VECTORS.
C SEED IS A REAL INPUT VARIABLE TO INITIALIZE THE RANDOM NUMBER GEN-
C ERATOR RANF USING THE INITIALIZING SUBROUTINE RANSET. IF SEED
C IDENTIFIES A NEGATIVE VALUE, NO PRINTING OF THE NON-ZERO
C DIAGONALS WILL BE DONE.
C TRANSA IS A REAL INPUT VARIABLE, THE LEFT ENDPOINT OF THE INTERVAL
C FROM WHICH THE NON-ZERO ENTRIES OF A AND T ARE SELECTED.
C DIALA IS A REAL INPUT VARIABLE, A POSITIVE NUMBER SELECTED SO THAT
C TRANSA + DIALA IS THE RIGHT ENDPOINT OF THE INTERVAL OVER

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ANMAX = 0.0 6630
L = 1 6640
220 DO 240 J = 1, MM 6650
    CALL BANDOP(N, NA, A, ND, X(1,J), WK(1,1)) 6660
    CALL TANDOP(N, NT, T, ND, X(1,J), WK(1,2)) 6670
    S = 0.0 6680
    DO 230 I = 1, N 6690
        S = S + (WK(I,1) - D(J)*WK(I,2))*(WK(I,1) - D(J)*WK(I,2)) 6700
230 CONTINUE 6710
    PROD = PROD*S 6720
    ANMAX = AMAX1(ANMAX, S) 6730
    L = MAX1(FLOAT(L), SIGN(FLOAT(J), S - ANMAX)) 6740
240 CONTINUE 6750
    S = ABS(D(1)) - D(IP) 6760
    ANMAX = SQRT(ANMAX)/S 6770
    ANMEAN = SQRT(PROD**(1.0/FLOAT(MM)))/S 6780
250 WRITE (6,260) IP, M, MM, KS, L, ANMAX, ANMEAN, T1 6790
260 FORMAT(1H0/1H0,57X,17HFINAL EIGENVALUES/1H /(1PE37.12,4E21.12)) 6800
270 CONTINUE 6810
280 CONTINUE 6820
    WRITE (6,290) (D(J), J = 1, MAXM) 6830
290 FORMAT(1H0/1H0,57X,17HFINAL EIGENVALUES/1H /(1PE37.12,4E21.12)) 6840
    GO TO 10 6850
    END 6860

SUBROUTINE INTROS(KS, G, H, F) 6870
C INTROS IS A DUMMY SUBROUTINE INF. 6880
    INTEGER G, H 6890
    REAL F(1) 6900
    RETURN 6910
    END 6920

SUBROUTINE BANDOP(N, NA, A, ND, V, W) 6930
C CALCULATE THE IMAGE W OF THE N-VECTOR V UNDER 6940
C THE MATRIX A WHERE A IS SYMMETRIC AND BANDED OF BAND WIDTH 2*NA + 1. 6950
    DIMENSION A(ND,1), V(1), W(1) 6960
    INTEGER P, Q, R 6970
    MP1 = NA + 1 6980
    DO 70 I = 1, N 6990
        P = MAX0(1, -I + (MP1 + 1)) 7000
        Q = MAX0(1, I + (MP1 - N)) 7010
        T = A(I,MP1)*V(I) 7020
        IF (NA - P) 30, 10, 10 7030
10 R = I - (MP1 - P) 7040
        DO 20 K = P, NA 7050
            T = T + A(I,K)*V(R) 7060
            R = R + 1 7070
20 CONTINUE 7080
30 IF (NA - Q) 60, 40, 40 7090
40 R = I + (MP1 - Q) 7100
        DO 50 K = Q, NA 7110
            T = T + A(R,K)*V(R) 7120
            R = R - 1 7130
50 CONTINUE 7140
60 W(I) = T 7150
70 CONTINUE 7160
    RETURN 7170
    END 7180

SUBROUTINE TANDOP(N, NA, A, ND, V, W) 7190
C COMPUTE THE IMAGE W OF THE N-VECTOR V UNDER THE N-SQUARE MATRIX 7200
C A(A-TRANPOSED), A LOWER TRIANGULAR OF BANDWIDTH NA + 1. 7210
    DIMENSION A(ND,1), V(1), W(1) 7220
    INTEGER P, Q, R 7230
    MP1 = NA + 1 7240
    DO 40 I = 1, N 7250
        P = MAX0(1, I + (MP1 - N)) 7260
        T = A(I,MP1)*V(I) 7270
        IF (NA - P) 30, 10, 10 7280
10 R = I + (MP1 - P) 7290
        DO 20 K = P, NA 7300
            T = T + A(R,K)*V(R) 7310
            R = R - 1 7320

```


20	CONTINUE	7330
30	W(I) = T	7340
40	CONTINUE	7350
	I = N	7360
	DO 80 Q = 1, N	7370
	P = MAX0(1, -I + (MP1 + 1))	7380
	T = A(I,MP1)*W(I)	7390
	IF (NA - P) 70, 50, 50	7400
50	R = I - (MP1 - P)	7410
	DO 60 K = P, NA	7420
	T = T + A(I,K)*W(R)	7430
	R = R + 1	7440
60	CONTINUE	7450
70	W(I) = T	7460
	I = I - 1	7470
80	CONTINUE	7480
	RETURN	7490
	END	7500
	FUNCTION BAND(N, Z, W)	7510
C	CALCULATE THE INNER PRODUCT (W-TRANPOSED)BZ WHERE B =	7520
C	T(T-TRANPOSED).	7530
	DIMENSION A(50,6), T(50,6), Y(50)	7540
	COMMON NT, T, NA, A, ND, Y	7550
	DIMENSION W(1), Z(1)	7560
10	CALL TANDOP(N, NT, T, ND, Z, Y)	7570
	S = 0.0	7580
	DO 15 I = 1, N	7590
	S = S + W(I)*Y(I)	7600
15	CONTINUE	7610
	BAND = S	7620
	RETURN	7630
	END	7640
	SUBROUTINE ABAND(N, Z, W)	7650
C	CALCULATE THE SOLUTION W OF THE LINEAR SYSTEM BW = AZ GIVEN THE	7660
C	N-VECTOR Z, THE BANDED SYMMETRIC N-SQUARE MATRIX A OF BAND WIDTH	7670
C	2*NA + 1, AND THE BANDED TRIANGULAR FACTOR T OF BAND WIDTH NT + 1	7680
C	OF THE POSITIVE DEFINITE MATRIX B, T(T-TRANPOSED) = B.	7690
	DIMENSION W(1), Z(1)	7700
	DIMENSION A(50,6), T(50,6), Y(50)	7710
	COMMON NT, T, NA, A, ND, Y	7720
	CALL BANDOP(N, NA, A, ND, Z, W)	7730
	CALL SOLVE(N, NT, T, ND, W)	7740
	RETURN	7750
	END	7760
	SUBROUTINE SOLVE(N, M, A, N1, B)	7770
C	SOLVE A LINEAR SYSTEM GIVEN A TRIANGULAR FACTORIZATION OF A BANDED	7780
C	POSITIVE DEFINITE COEFFICIENT MATRIX OF BAND WIDTH 2*M + 1.	7790
	DIMENSION A(N1,1), B(N)	7800
	INTEGER P, Q, R	7810
	MP1 = M + 1	7820
COMMENT	SOLUTION OF LY = B	7830
	DO 130 I = 1, N	7840
	P = MAX0(1, MP1 - I + 1)	7850
	T = B(I)	7860
	IF (M - P) 120, 100, 100	7870
100	Q = I - MP1 + P	7880
	DO 110 K = P, M	7890
	T = T - A(I,K)*B(Q)	7900
	Q = Q + 1	7910
110	CONTINUE	7920
120	B(I) = T/A(I,MP1)	7930
130	CONTINUE	7940
COMMENT	SOLUTION OF UX = Y	7950
	I = N	7960
	DO 170 R = 1, N	7970
	P = MAX0(1, MP1 - N + I)	7980
	T = B(I)	7990
	IF (M - P) 160, 140, 140	8000
140	Q = I + MP1 - P	8010
	DO 150 K = P, M	8020

```

      T = T - A(Q,K)*B(Q)
      Q = Q - 1
150  CONTINUE
160  B(I) = T/A(I,MP1)
      I = I - 1
170  CONTINUE
      RETURN
      END

```

8030
8040
8050
8060
8070
8080
8090
8100

```

      FUNCTION RANF(X)
C RETURNS A RANDOM NUMBER ON THE OPEN UNIT INTERVAL GIVEN ANY ARGUMENT.
      COMMON /TSEED/ K
      K = MOD(3125*K, 65536)
      RANF = ABS(FLOAT(K - 32768)/32768.0)
      RETURN
      END

```

8110
8120
8130
8140
8150
8160
8170

```

      SUBROUTINE RANSET(SEED)
C INITIALIZES THE RANDOM NUMBER SEED K USING THE REAL INPUT
C VARIABLE SEED.
      COMMON /TSEED/ K
      T = SEED
      IF (ABS(SEED) .GT. 1.0) T = 1.0/T
      K = 2*IFIX(16384.0*T) - IFIX(SIGN(1.0, T))
      RETURN
      END

```

8180
8190
8200
8210
8220
8230
8240
8250
8260

```

      FUNCTION SECOND(T)
C DUMMY FUNCTION SECOND RETURNS THE VALUES SECOND = T = 0.0.
C
C SECOND MAY BE MODIFIED TO CALL AN OPERATING SYSTEM DEPENDENT PROGRAM
C WHICH RETURNS CENTRAL PROCESSOR TIME IN SECONDS RELATIVE TO
C AN ARBITRARY STARTING POINT AS A REAL VALUE T.
C
      T = 0.0
      SECOND = 0.0
      RETURN
      END

```

8270
8280
8290
8300
8310
8320
8330
8340
8350
8360
8370

```

C
C -----
C
C SUBROUTINE IMTQL2(NM,N,D,E,Z,IERR)
C
C INTEGER I,J,K,L,M,N,II,NM,MML,IERR
C REAL D(N),E(N),Z(NM,N)
C REAL B,C,F,G,P,R,S,MACHEP
C REAL SQRT,ABS,SIGN
C
C THIS SUBROUTINE IS A TRANSLATION OF THE ALGOL PROCEDURE IMTQL2,
C NUM. MATH. 12, 377-383(1968) BY MARTIN AND WILKINSON,
C AS MODIFIED IN NUM. MATH. 15, 450(1970) BY DUBRULLE.
C HANDBOOK FOR AUTO. COMP., VOL.II-LINEAR ALGEBRA, 241-248(1971).
C
C THIS SUBROUTINE FINDS THE EIGENVALUES AND EIGENVECTORS
C OF A SYMMETRIC TRIDIAGONAL MATRIX BY THE IMPLICIT QL METHOD.
C THE EIGENVECTORS OF A FULL SYMMETRIC MATRIX CAN ALSO
C BE FOUND IF TRED2 HAS BEEN USED TO REDUCE THIS
C FULL MATRIX TO TRIDIAGONAL FORM.
C
C ON INPUT-
C
C NM MUST BE SET TO THE ROW DIMENSION OF TWO-DIMENSIONAL
C ARRAY PARAMETERS AS DECLARED IN THE CALLING PROGRAM
C DIMENSION STATEMENT,
C
C N IS THE ORDER OF THE MATRIX,
C
C D CONTAINS THE DIAGONAL ELEMENTS OF THE INPUT MATRIX,
C
C E CONTAINS THE SUBDIAGONAL ELEMENTS OF THE INPUT MATRIX

```

8380
8390
8400
8410
8420
8430
8440
8450
8460
8470
8480
8490
8500
8510
8520
8530
8540
8550
8560
8570
8580
8590
8600
8610
8620
8630
8640
8650
8660
8670
8680
8690

```

C           IN ITS LAST N-1 POSITIONS. E(1) IS ARBITRARY,      8700
C
C           Z CONTAINS THE TRANSFORMATION MATRIX PRODUCED IN THE  8710
C           REDUCTION BY TRED2, IF PERFORMED. IF THE EIGENVECTORS  8720
C           OF THE TRIDIAGONAL MATRIX ARE DESIRED, Z MUST CONTAIN  8730
C           THE IDENTITY MATRIX.                                8740
C
C           ON OUTPUT-                                         8750
C
C           D CONTAINS THE EIGENVALUES IN ASCENDING ORDER. IF AN  8760
C           ERROR EXIT IS MADE, THE EIGENVALUES ARE CORRECT BUT  8770
C           UNORDERED FOR INDICES 1,2,...,IERR-1,              8780
C
C           E HAS BEEN DESTROYED,                               8790
C
C           Z CONTAINS ORTHONORMAL EIGENVECTORS OF THE SYMMETRIC  8800
C           TRIDIAGONAL (OR FULL) MATRIX. IF AN ERROR EXIT IS MADE,  8810
C           Z CONTAINS THE EIGENVECTORS ASSOCIATED WITH THE STORED  8820
C           EIGENVALUES,                                       8830
C
C           IERR IS SET TO                                     8840
C           ZERO FOR NORMAL RETURN,                             8850
C           J IF THE J-TH EIGENVALUE HAS NOT BEEN              8860
C           DETERMINED AFTER 30 ITERATIONS.                     8870
C
C           QUESTIONS AND COMMENTS SHOULD BE DIRECTED TO B. S. GARBOW,  8880
C           APPLIED MATHEMATICS DIVISION, ARGONNE NATIONAL LABORATORY  8890
C
C           -----                                           8900
C           ***** MACHEP IS A MACHINE DEPENDENT PARAMETER SPECIFYING  8910
C           THE RELATIVE PRECISION OF FLOATING POINT ARITHMETIC.  8920
C
C           *****                                           8930
C           MACHEP = 2.**(-47)                                    8940
C
C           IERR = 0                                           8950
C           IF (N .EQ. 1) GO TO 1001                             8960
C
C           DO 100 I = 2, N                                       8970
1000 E(I-1) = E(I)                                             8980
C
C           E(N) = 0.0                                          8990
C
C           DO 240 L = 1, N                                       9000
C           J = 0                                               9010
C           ***** LOOK FOR SMALL SUB-DIAGONAL ELEMENT *****  9020
1005 DO 110 M = L, N                                           9030
C           IF (M .EQ. N) GO TO 120                               9040
C           IF (ABS(E(M)) .LE. MACHEP * (ABS(D(M)) + ABS(D(M+1))))  9050
X           GO TO 120                                           9060
1100 CONTINUE                                                 9070
C
C           120 P = D(L)                                         9080
C           IF (M .EQ. L) GO TO 240                               9090
C           IF (J .EQ. 30) GO TO 10000                          9100
C           J = J + 1                                           9110
C
C           ***** FORM SHIFT *****                           9120
C           G = (D(L+1) - P) / (2.0 * E(L))                      9130
C           R = SQRT(G*G+1.0)                                     9140
C           G = D(M) - P + E(L) / (G + SIGN(R,G))                9150
C           S = 1.0                                             9160
C           C = 1.0                                             9170
C           P = 0.0                                             9180
C           MML = M - L                                         9190
C
C           ***** FOR I=M-1 STEP -1 UNTIL L DO -- *****  9200
C           DO 200 II = 1, MML                                    9210
C           I = M - II                                          9220
C           F = S * E(I)                                         9230
C           B = C * E(I)                                         9240
C           IF (ABS(F) .LT. ABS(G)) GO TO 150                    9250
C           C = G / F                                           9260
C           R = SQRT(C*C+1.0)                                     9270
C           E(I+1) = F * R                                       9280
C           S = 1.0 / R                                         9290
C           C = C * S                                           9300

```

```

          GO TO 160
150      S = F / G
          R = SQRT(S*S+1.0)
          E(I+1) = G * R
          C = 1.0 / R
          S = S * C
160      G = D(I+1) - P
          R = (D(I) - G) * S + 2.0 * C * B
          P = S * R
          D(I+1) = G + P
          G = C * R - B
C      ***** FORM VECTOR *****
          DO 180 K = 1, N
            F = Z(K,I+1)
            Z(K,I+1) = S * Z(K,I) + C * F
            Z(K,I) = C * Z(K,I) - S * F
180      CONTINUE
C
200      CONTINUE
C
          D(L) = D(L) - P
          E(L) = G
          E(M) = 0.0
          GO TO 105
240      CONTINUE
C      ***** ORDER EIGENVALUES AND EIGENVECTORS *****
          DO 300 II = 2, N
            I = II - 1
            K = I
            P = D(I)
C
          DO 260 J = II, N
            IF (D(J) .GE. P) GO TO 260
            K = J
            P = D(J)
260      CONTINUE
C
          IF (K .EQ. I) GO TO 300
          D(K) = D(I)
          D(I) = P
C
          DO 280 J = 1, N
            P = Z(J,I)
            Z(J,I) = Z(J,K)
            Z(J,K) = P
280      CONTINUE
C
300      CONTINUE
C
          GO TO 1001
C      ***** SET ERROR -- NO CONVERGENCE TO AN
C      EIGENVALUE AFTER 30 ITERATIONS *****
1000     IERR = L
1001     RETURN
C      ***** LAST CARD OF IMTQL2 *****
          END
C
C      -----
C
          SUBROUTINE TRED2(NM,N,A,D,E,Z)
C
          INTEGER I,J,K,L,N,II,NM,JP1
          REAL A(NM,N),D(N),E(N),Z(NM,N)
          REAL F,G,H,HH,SCALE
          REAL SQRT,ABS,SIGN
C
C      THIS SUBROUTINE IS A TRANSLATION OF THE ALGOL PROCEDURE TRED2,
C      NUM. MATH. 11, 181-195(1968) BY MARTIN, REINSCH, AND WILKINSON.
C      HANDBOOK FOR AUTO. COMP., VOL.II-LINEAR ALGEBRA, 212-226(1971).
C
C      THIS SUBROUTINE REDUCES A REAL SYMMETRIC MATRIX TO A
C      SYMMETRIC TRIDIAGONAL MATRIX USING AND ACCUMULATING
C      ORTHOGONAL SIMILARITY TRANSFORMATIONS.
C
          ON INPUT-

```

```

9460
9470
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9490
9500
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9520
9530
9540
9550
9560
9570
9580
9590
9600
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9620
9630
9640
9650
9660
9670
9680
9690
9700
9710
9720
9730
9740
9750
9760
9770
9780
9790
9800
9810
9820
9830
9840
9850
9860
9870
9880
9890
9900
9910
9920
9930
9940
9950
9960
9970
9980
9990
10000
10010
10020
10030
10040
10050
10060
10070
10080
10090
10100
10110
10120
10130
10140
10150
10160
10170
10180
10190
10200

```

```

C
C      NM MUST BE SET TO THE ROW DIMENSION OF TWO-DIMENSIONAL
C      ARRAY PARAMETERS AS DECLARED IN THE CALLING PROGRAM
C      DIMENSION STATEMENT,
C
C      N IS THE ORDER OF THE MATRIX,
C
C      A CONTAINS THE REAL SYMMETRIC INPUT MATRIX. ONLY THE
C      LOWER TRIANGLE OF THE MATRIX NEED BE SUPPLIED.
C
C      ON OUTPUT-
C
C      D CONTAINS THE DIAGONAL ELEMENTS OF THE TRIDIAGONAL MATRIX,
C
C      E CONTAINS THE SUBDIAGONAL ELEMENTS OF THE TRIDIAGONAL
C      MATRIX IN ITS LAST N-1 POSITIONS. E(1) IS SET TO ZERO,
C
C      Z CONTAINS THE ORTHOGONAL TRANSFORMATION MATRIX
C      PRODUCED IN THE REDUCTION,
C
C      A AND Z MAY COINCIDE. IF DISTINCT, A IS UNALTERED.
C
C      QUESTIONS AND COMMENTS SHOULD BE DIRECTED TO B. S. GARBOW,
C      APPLIED MATHEMATICS DIVISION, ARGONNE NATIONAL LABORATORY
C
C      -----
C
C      DO 100 I = 1, N
C
C          DO 100 J = 1, I
C              Z(I,J) = A(I,J)
C
C      100 CONTINUE
C
C      IF (N .EQ. 1) GO TO 320
C      ***** FOR I=N STEP -1 UNTIL 2 DO -- *****
C      DO 300 II = 2, N
C          I = N + 2 - II
C          L = I - 1
C          H = 0.0
C          SCALE = 0.0
C          IF (L .LT. 2) GO TO 130
C      ***** SCALE ROW (ALGOL TOL THEN NOT NEEDED) *****
C      DO 120 K = 1, L
C      120  SCALE = SCALE + ABS(Z(I,K))
C
C          IF (SCALE .NE. 0.0) GO TO 140
C      130  E(I) = Z(I,L)
C          GO TO 290
C
C      140  DO 150 K = 1, L
C              Z(I,K) = Z(I,K) / SCALE
C              H = H + Z(I,K) * Z(I,K)
C      150  CONTINUE
C
C          F = Z(I,L)
C          G = -SIGN(SQRT(H),F)
C          E(I) = SCALE * G
C          H = H - F * G
C          Z(I,L) = F - G
C          F = 0.0
C
C          DO 240 J = 1, L
C              Z(J,I) = Z(I,J) / H
C              G = 0.0
C      ***** FORM ELEMENT OF A*U *****
C      DO 180 K = 1, J
C      180  G = G + Z(J,K) * Z(I,K)
C
C          JP1 = J + 1
C          IF (L .LT. JP1) GO TO 220
C
C          DO 200 K = JP1, L
C      200  G = G + Z(K,J) * Z(I,K)
C      ***** FORM ELEMENT OF P *****
C      220  E(J) = G / H
C          F = F + E(J) * Z(I,J)

```

10210
10220
10230
10240
10250
10260
10270
10280
10290
10300
10310
10320
10330
10340
10350
10360
10370
10380
10390
10400
10410
10420
10430
10440
10450
10460
10470
10480
10490
10500
10510
10520
10530
10540
10550
10560
10570
10580
10590
10600
10610
10620
10630
10640
10650
10660
10670
10680
10690
10700
10710
10720
10730
10740
10750
10760
10770
10780
10790
10800
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10820
10830
10840
10850
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10870
10880
10890
10900
10910
10920
10930
10940
10950
10960

```

240 CONTINUE
C
      HH = F / (H + H)
C ***** FORM REDUCED A *****
      DO 260 J = 1, L
        F = Z(I,J)
        G = E(J) - HH * F
        E(J) = G
C
      DO 260 K = 1, J
        Z(J,K) = Z(J,K) - F * E(K) - G * Z(I,K)
260 CONTINUE
C
290 D(I) = H
300 CONTINUE
C
320 D(1) = 0.0
      E(1) = 0.0
C ***** ACCUMULATION OF TRANSFORMATION MATRICES *****
      DO 500 I = 1, N
        L = I - 1
        IF (D(I) .EQ. 0.0) GO TO 380
C
        DO 360 J = 1, L
          G = 0.0
C
          DO 340 K = 1, L
            G = G + Z(I,K) * Z(K,J)
C
          DO 360 K = 1, L
            Z(K,J) = Z(K,J) - G * Z(K,I)
360 CONTINUE
C
380 D(I) = Z(I,I)
      Z(I,I) = 1.0
      IF (L .LT. 1) GO TO 500
C
      DO 400 J = 1, L
        Z(I,J) = 0.0
        Z(J,I) = 0.0
400 CONTINUE
C
500 CONTINUE
C
      RETURN
C ***** LAST CARD OF TRED2 *****
      END

```

30	5	4	200	1.00E-10	1.00E-01	-1.00E+01	2.00E+01		DATA
30	5	4	-200	1.00E-10	-1.00E-01	-1.00E+01	2.00E+01		DATA
40	4	3	200	1.00E-10	2.00E-01	-1.00E+02	2.00E+02		DATA
40	4	3	-200	1.00E-10	-2.00E-01	-1.00E+02	2.00E+02		DATA
50	2	1	200	1.00E-10	3.00E-01	-1.00E+02	2.00E+02		DATA
50	2	1	-200	1.00E-10	-3.00E-01	-1.00E+02	2.00E+02		DATA

1	N	NA	NB	KM	EPS
	30	5	4	200	1.00E-10

ADJACENT NON-ZERO LOWER DIAGONALS

	7	7	7	5	10	9	10	6	9
	5	9	7	7	5	8	7	6	8
	8	5	-5	7	-5	3	-4	8	-9
	5	6	-4	-5	-9	-7	2	5	-3
	4	5	3	3	-6	9	-6	-6	2
	0	-2	2	7	6	-1	4	6	-1
	0	-7	-2	0	-9	4	4	-9	8
	-8	-7	0	4	6	-8	-5	-3	-5
	-6	9	6	-4	3	7	2	1	-1
	1	-7	-8	8	9	-5	3	3	5

ADJACENT NON-ZERO LOWER DIAGONALS

∅	7	6	-4	-6	5	8	3	-7
-8	1	-4	6	-9	8	5	3	-4
-7	-1	∅	4	8	1	-5	∅	5
-4	∅	4	4	5	-2	2	-7	-3
-2	-1	-8	7	1	∅	-6	6	9
2	1	9	5	-6	-6	-6	∅	-9
∅	∅	-8	∅	-9	∅	∅	7	6
6	5	-6	9	-8	5	-4	-9	6
-7	∅	-5	8	∅	-6	-9	5	-7
6	-5	2	∅	6	-4	-5	-2	-4
-8	∅	8	3	9	4	1	9	-8
3	3	-8	∅	2	-6	-4	∅	4

∅
∅

P	EM(IN)	EM(OUT)	KS	K	MAX RES
2	1	1	29	1	8.6
3	1	1	15	1	6.4
3	2	2	24	1	4.9
4	1	1	15	1	1.1
4	2	2	15	1	1.3
4	3	3	51	1	1.2
5	1	1	13	1	9.∅
5	2	2	13	1	1.1
5	3	3	21	1	6.2
5	4	4	21	1	1.3
6	1	1	13	1	1.7
6	2	2	13	1	5.4
6	3	3	21	1	7.2
6	4	4	13	1	9.6
6	5	5	4∅	1	1.2

∅
∅

FINAL EIGENVALUES

1	N	4.55∅1∅9524972E+∅1	2.4271∅357641∅E+∅1	-7.628352951∅42E+∅∅
		NA	NB	EPS
	3∅	5	4	-2∅∅
				1.∅∅E-1∅

∅
∅
∅

P	EM(IN)	EM(OUT)	KS	K	MAX RES
2	1	1	3∅	1	8.4
3	1	1	3	1	5.7
3	2	2	15	1	1.1
4	1	1	3	1	1.2
4	2	2	3	1	9.3
4	3	3	31	1	8.2
5	1	1	3	1	7.3
5	2	2	3	1	5.∅
5	3	3	3	1	9.1
5	4	4	3	1	6.3
6	1	1	3	1	6.2
6	2	2	3	1	1.3
6	3	3	3	1	8.5
6	4	4	3	1	1.1
6	5	5	18	1	1.3

∅
∅

FINAL EIGENVALUES

1	N	4.55∅1∅9524972E+∅1	2.4271∅357641∅E+∅1	-7.628352951∅42E+∅∅
		NA	NB	EPS
	4∅	4	3	2∅∅
				1.∅∅E-1∅

∅
∅

ADJACENT NON-ZERO LOWER DIAGONALS

5∅	62	67	76	78	67	67	98	1∅∅
82	66	64	84	66	89	76	5∅	6∅
58	61	99	56	1∅∅	91	55	83	68

72	-73	-61	70	-94	-36	-80	29	-13
-3	76	25	-85	78	21	37	3	77
-15	52	-96	76	-91	26	52	-39	72
50	16	35	20	73	-50	76	-75	-27
-89	95	-96	60	90	-81	-71	-71	-45
-19	80	40	23	-44	39	35	-91	
-76	97	-47	47	36	74	-44	-18	-71
24	34	-61	-86	66	-25	10	16	32
-3	10	55	25	-74	-19	14		

0

ADJACENT NON-ZERO LOWER DIAGONALS

-53	73	-15	-63	70	39	12	-48	-99
13	-21	78	-34	-17	62	64	-17	-64
85	-13	6	12	-20	-76	-20	-57	77
66	77	-46	66	-49	-86	40	-14	69
8	-93	-93	-98	28	51	40	-33	53
73	27	-14	96	-31	-15	-71	6	17
64	2	49	46	-60	24	63	-50	-22
54	99	-5	51	-9	-43	-34	-46	-98
15	-7	-2	33	77	41	72	-81	
-66	-40	42	-45	10	-97	-57	-14	-67
-48	12	93	0	21	-79	7	3	-54
62	-5	16	-52	82	-28	43		
-77	-79	-4	-63	61	31	-3	55	-24
83	8	-44	81	-68	-44	9	-8	-68
-64	8	19	-35	42	99			

0
0

P	EM(IN)	EM(OUT)	KS	K	MAX RES
2	1	1	13	1	2.8
3	1	1	13	1	4.2
3	2	2	13	1	4.3
4	1	1	7	1	5.0
4	2	2	13	1	4.1
4	3	3	54	1	3.8
5	1	1	7	1	4.9
5	2	2	13	1	5.2
5	3	3	34	1	4.0
5	4	4	170	1	5.2
6	1	1	7	1	4.8
6	2	2	13	1	4.4
6	3	3	30	1	4.6
6	4	4	37	1	5.0
6	5	5	26	1	6.7
7	1	1	7	1	5.6
7	2	2	13	1	5.6
7	3	3	25	1	5.9
7	4	4	27	1	5.0
7	5	5	27	1	5.0
7	6	5	91	1	6.6
8	1	1	7	1	5.7
8	2	2	13	1	5.5
8	3	3	30	1	6.9
8	4	4	37	1	6.1
8	5	5	28	1	7.8
8	6	6	180	1	6.6
8	7	6	177	1	7.8

0
0

FINAL EIGENVALUES

-1.059611108016E+03	-2.114191261132E+01	2.825002039426E-01
-6.809394216930E-02	-6.243506342913E-02	

1 N
40

NA	NB	KM	EPS		
4	3	-200	1.00E-10		
P	EM(IN)	EM(OUT)	KS	K	MAX RES
2	1	1	13	1	3.1

0

3	1	1	3	1	3.4
3	2	2	3	1	3.8
4	1	1	3	1	5.6
4	2	2	7	1	4.1
4	3	3	40	1	4.4
5	1	1	3	1	4.0
5	2	2	3	1	3.6
5	3	3	18	1	5.2
5	4	4	100	1	3.8
6	1	1	3	1	6.2
6	2	2	3	1	4.9
6	3	3	11	1	2.8
6	4	4	13	1	5.1
6	5	5	13	1	4.7
7	1	1	3	1	3.9
7	2	2	3	1	4.0
7	3	3	7	1	5.7
7	4	4	19	1	3.8
7	5	5	13	1	5.4
7	6	6	187	1	4.2
8	1	1	3	1	4.1
8	2	2	3	1	5.9
8	3	3	7	1	5.5
8	4	4	10	1	4.7
8	5	5	12	1	6.1
8	6	6	55	1	6.1
8	7	6	157	1	4.9

0
0

FINAL EIGENVALUES

-1.059611108016E+03	-2.114191261132E+01	2.825002039426E-01
-6.809394215568E-02	-6.243506344448E-02	

1

N

NA

NB

KM

EPS

50

2

1

200

1.00E-10

0
0

ADJACENT NON-ZERO LOWER DIAGONALS

74	79	100	89	95	93	60	59	57
71	61	74	78	70	95	78	64	76
51	96	98	71	55	72	84	51	66
72	95	62	68	82				
46	11	-81	16	-40	50	-96	12	59
59	90	-54	17	88	-97	-18	79	18
-20	-20	73	2	-16	94	64	86	44
92	-5	1	-63					

0

ADJACENT NON-ZERO LOWER DIAGONALS

-40	84	-74	54	-91	80	-90	-60	-92
90	74	-97	-46	-43	-95	37	-12	-89
61	-57	-55	13	44	41	-57	-79	-91
-27	77	-98	62	49				
-67	0	-19	68	-12	50	53	-44	-56
-6	-63	-4	-86	-42	2	-26	-3	-50
-38	73	35	-83	32	-58	-32	28	-20
73	26	-66	-18					
13	-74	-67	-20	-77	53	92	-74	71
57	61	56	-91	-62	-90	-14	45	-90
62	30	-37	88	25	-70	-33	-99	26
-34	-47	-78						

0
0

P	EM(IN)	EM(OUT)	KS	K	MAX RES
2	1	1	46	1	7.0
3	1	1	43	1	1.8
3	2	2	60	2	2.7
4	1	1	30	1	7.0
4	2	2	43	1	9.9
4	3	3	63	3	2.5
5	1	1	26	1	1.1
5	2	2	39	2	2.6
5	3	3	52	1	2.3
5	4	4	61	1	3.0

6	1	1	26	1	1.4
6	2	2	39	1	6.5
6	3	3	39	3	5.7
6	4	4	50	2	1.7
6	5	5	107	1	7.7
7	1	1	24	1	1.2
7	2	2	24	1	1.2
7	3	3	24	3	4.4
7	4	4	33	1	1.3
7	5	5	35	3	3.9
7	6	6	35	3	1.4
8	1	1	24	1	1.4
8	2	2	24	1	1.6
8	3	3	24	1	1.4
8	4	4	24	4	3.9
8	5	5	35	4	2.6
8	6	6	35	4	1.8
8	7	7	114	4	3.2
9	1	1	24	1	1.4
9	2	2	24	1	2.2
9	3	3	24	1	2.3
9	4	4	24	1	2.2
9	5	5	35	4	1.4
9	6	6	24	6	5.0
9	7	7	63	5	1.5
9	8	7	63	7	4.2
10	1	1	24	1	1.4
10	2	2	24	1	1.5
10	3	3	24	1	1.8
10	4	4	24	1	1.5
10	5	5	24	1	1.5
10	6	6	24	6	1.7
10	7	7	48	1	2.8
10	8	8	48	8	4.2
10	9	9	48	9	1.0

0
0

FINAL EIGENVALUES

-2.536352449033E-01	2.038265111579E-01	-1.695641085298E-01
-1.194676112328E-01	7.012023880100E-02	6.100076826561E-02
NA 2	KM 1 -200	EPS 10

1
0
0
0
0
N
50

P	EM(IN)	EM(OUT)	KS	K	MAX RES
2	1	1	44	1	2.0
3	1	1	3	1	1.8
3	2	2	43	2	6.5
4	1	1	3	1	7.9
4	2	2	3	1	1.9
4	3	3	56	3	1.5
5	1	1	3	1	9.6
5	2	2	3	1	1.8
5	3	3	3	1	1.7
5	4	4	44	4	3.5
6	1	1	3	1	1.0
6	2	2	3	1	1.8
6	3	3	3	1	1.5
6	4	4	3	1	1.8
6	5	5	107	5	2.2
7	1	1	3	1	1.0
7	2	2	3	1	1.2
7	3	3	3	1	1.7
7	4	4	3	1	1.7
7	5	5	8	1	1.7
7	6	6	6	6	1.6
8	1	1	3	1	1.7
8	2	2	3	1	1.8
8	3	3	3	1	1.5
8	4	4	3	1	2.0
8	5	5	4	1	2.1
8	6	6	3	1	1.6
8	7	7	80	7	1.8
9	1	1	3	1	1.5
9	2	2	3	1	1.8
9	3	3	3	1	1.9

COLLECTED ALGORITHMS (cont.)

538-P25- 0

9	4	4	3	1	1.4
9	5	5	3	1	1.8
9	6	6	3	1	1.3
9	7	7	3	1	1.7
9	8	8	99	1	1.7
10	1	1	3	1	1.6
10	2	2	3	1	1.4
10	3	3	3	1	1.5
10	4	4	3	1	1.7
10	5	5	3	1	1.6
10	6	6	3	1	1.7
10	7	7	3	1	1.8
10	8	8	18	1	1.8
10	9	9	3	1	1.6

0
0

FINAL EIGENVALUES

-2.536352449033E-01	2.038265111579E-01	-1.695641085298E-01
-1.194676112328E-01	7.012023880101E-02	6.100076826562E-02

ALGORITHM 539

Basic Linear Algebra Subprograms for Fortran Usage [F1]

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Key Words and Phrases: linear algebra, utilities

CR Categories: 4.49, 5.14

Language: Fortran, assembly language

DESCRIPTION

This package complements [1], where further details are given.

REFERENCES

1. LAWSON, C.L., HANSON, R.J., KINCAID, D.R., AND KROGH, F.T. Basic linear algebra subprograms for Fortran usage. *ACM Trans. Math. Software* 5, 3 (September 1979), 308-323.

ALGORITHM

[The contents of this package consists of four files. The first file contains Fortran versions for the BLAS (38 subprograms), programs for testing the BLAS (13 modules), and 18 subprograms from Brent's multiple precision package that are used in the implementation of the extended precision inner products. The remaining 3 files contain Fortran callable assembly language versions for 3 different machines: IBM 360/370 series, CDC 6000 series, and Univac 1100 series. Printed here are the 38 BLAS subprograms. The complete listing is available from the ACM Algorithms Distribution Service.]

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```

REAL FUNCTION SDOT(N,SX,INCX,SY,INCY)
C
C RETURNS THE DOT PRODUCT OF SINGLE PRECISION SX AND SY.
C SDOT = SUM FOR I = 0 TO N-1 OF SX(LX+I*INCX) * SY(LY+I*INCY),
C WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N, AND LY IS
C DEFINED IN A SIMILAR WAY USING INCY.
C
REAL SX(1),SY(1)
SDOT = 0.0E0
IF(N.LE.0)RETURN
IF(INCX.EQ.INCY) IF(INCX-1)5,20,60
5 CONTINUE
C
C CODE FOR UNEQUAL INCREMENTS OR NONPOSITIVE INCREMENTS.
C
IX = 1
IY = 1
IF(INCX.LT.0)IX = (-N+1)*INCX + 1
IF(INCY.LT.0)IY = (-N+1)*INCY + 1
DO 10 I = 1,N
SDOT = SDOT + SX(IX)*SY(IY)
IX = IX + INCX
IY = IY + INCY
10 CONTINUE
RETURN
C
C CODE FOR BOTH INCREMENTS EQUAL TO 1
C
C CLEAN-UP LOOP SO REMAINING VECTOR LENGTH IS A MULTIPLE OF 5.
C
20 M = MOD(N,5)
IF( M .EQ. 0 ) GO TO 40
DO 30 I = 1,M
SDOT = SDOT + SX(I)*SY(I)
30 CONTINUE
IF( N .LT. 5 ) RETURN
40 MP1 = M + 1
DO 50 I = MP1,N,5
SDOT = SDOT + SX(I)*SY(I) + SX(I + 1)*SY(I + 1) +
$ SX(I + 2)*SY(I + 2) + SX(I + 3)*SY(I + 3) + SX(I + 4)*SY(I + 4)
50 CONTINUE
RETURN
C
C CODE FOR POSITIVE EQUAL INCREMENTS .NE.1.
C
60 CONTINUE
NS=N*INCX
DO 70 I=1,NS,INCX
SDOT = SDOT + SX(I)*SY(I)
70 CONTINUE
RETURN
END

```

DOUBLE PRECISION FUNCTION DSDOT(N,SX,INCX,SY,INCY)

```

C
C RETURNS D.P. DOT PRODUCT ACCUMULATED IN D.P., FOR S.P. SX AND SY
C DSDOT = SUM FOR I = 0 TO N-1 OF SX(LX+I*INCX) * SY(LY+I*INCY),
C WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N, AND LY IS
C DEFINED IN A SIMILAR WAY USING INCY.
C
REAL SX(1),SY(1)
C
DSDOT = 0.D0
IF(N .LE. 0)RETURN
IF(INCX.EQ.INCY.AND.INCX.GT.0) GO TO 20
KX = 1
KY = 1
IF(INCX.LT.0) KX = 1+(1-N)*INCX
IF(INCY.LT.0) KY = 1+(1-N)*INCY
DO 10 I = 1,N
DSDOT = DSDOT + DBLE(SX(KX))*DBLE(SY(KY))
KX = KX + INCX
KY = KY + INCY
20

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```

10 CONTINUE                                DS18100
   RETURN                                  DS18110
20 CONTINUE                                DS18120
   NS = N*INCX                             DS18130
     DO 30 I=1,NS,INCX                     DS18140
       DSDOT = DSDOT + DBLE(SX(I))*DBLE(SY(I)) DS18150
30 CONTINUE                                DS18160
   RETURN                                  DS18170
   END                                     DS18180

REAL FUNCTION SDSDOT(N,SB,SX,INCX,SY,INCY) SD18190
C
C RETURNS S.P. RESULT WITH DOT PRODUCT ACCUMULATED IN D.P.
C SDSDOT = SB + SUM FOR I = 0 TO N-1 OF SX(LX+I*INCX)*SY(LY+I*INCY),
C WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N, AND LY IS
C DEFINED IN A SIMILAR WAY USING INCY.
C
REAL          SX(1),SY(1),SB              SD18240
DOUBLE PRECISION DSDOT                   SD18250
C
DSDOT = DBLE(SB)                         SD18260
IF(N .LE. 0) GO TO 30                    SD18270
IF(INCX.EQ.INCY.AND.INCX.GT.0) GO TO 40 SD18280
KX = 1                                    SD18290
KY = 1                                    SD18300
IF(INCX.LT.0) KX = 1+(1-N)*INCX          SD18310
IF(INCY.LT.0) KY = 1+(1-N)*INCY          SD18320
DO 10 I = 1,N                             SD18330
  DSDOT = DSDOT + DBLE(SX(KX))*DBLE(SY(KY)) SD18340
  KX = KX + INCX                           SD18350
  KY = KY + INCY                           SD18360
10 CONTINUE                               SD18370
30 SDSDOT = SNGL(DSDOT)                   SD18380
   RETURN                                  SD18390
40 CONTINUE                               SD18400
   NS = N*INCX                             SD18410
     DO 50 I=1,NS,INCX                     SD18420
       DSDOT = DSDOT + DBLE(SX(I))*DBLE(SY(I)) SD18430
50 CONTINUE                               SD18440
   SDSDOT = SNGL(DSDOT)                   SD18450
   RETURN                                  SD18460
   END                                     SD18470
                                         SD18480

DOUBLE PRECISION FUNCTION DDOT(N,DX,INCX,DY,INCY) DD18490
C
C RETURNS THE DOT PRODUCT OF DOUBLE PRECISION DX AND DY.
C DDOT = SUM FOR I = 0 TO N-1 OF DX(LX+I*INCX) * DY(LY+I*INCY)
C WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N, AND LY IS
C DEFINED IN A SIMILAR WAY USING INCY.
C
C
C DOUBLE PRECISION DX(1),DY(1)            DD18520
DDOT = 0.D0                               DD18530
IF(N.LE.0) RETURN                         DD18540
IF(INCX.EQ.INCY) IF(INCX-1) 5,20,60      DD18550
5 CONTINUE                                DD18560
                                         DD18570
C
C CODE FOR UNEQUAL OR NONPOSITIVE INCREMENTS.
C
C IX = 1                                   DD18580
C IY = 1                                   DD18590
C IF(INCX.LT.0) IX = (-N+1)*INCX + 1      DD18600
C IF(INCY.LT.0) IY = (-N+1)*INCY + 1      DD18610
C DO 10 I = 1,N                             DD18620
  DDOT = DDOT + DX(IX)*DY(IY)             DD18630
  IX = IX + INCX                           DD18640
  IY = IY + INCY                           DD18650
10 CONTINUE                               DD18660
   RETURN                                  DD18670
C
C CODE FOR BOTH INCREMENTS EQUAL TO 1.
C
C CLEAN-UP LOOP SO REMAINING VECTOR LENGTH IS A MULTIPLE OF 5. DD18680
                                         DD18690
                                         DD18700
                                         DD18710
                                         DD18720
                                         DD18730
                                         DD18740
                                         DD18750

```

```

C
20 M = MOD(N,5)
IF( M .EQ. 0 ) GO TO 40
DO 30 I = 1,M
  DDOT = DDOT + DX(I)*DY(I)
30 CONTINUE
IF( N .LT. 5 ) RETURN
40 MP1 = M + 1
DO 50 I = MP1,N,5
  DDOT = DDOT + DX(I)*DY(I) + DX(I+1)*DY(I+1) +
  $ DX(I + 2)*DY(I + 2) + DX(I + 3)*DY(I + 3) + DX(I + 4)*DY(I + 4)
50 CONTINUE
RETURN

C
C      CODE FOR POSITIVE EQUAL INCREMENTS .NE.1.
C
60 CONTINUE
NS = N*INCX
DO 70 I=1,NC,INCY
  DDOT = DDOT + DX(I)*DY(I)
70 CONTINUE
RETURN
END

DOUBLE PRECISION FUNCTION DQDOTA(N,DB,QC,DX,INCX,DY,INCY)
C
C D.P. DOT PRODUCT WITH EXTENDED PRECISION ACCUMULATION (AND RESULT)
C
C QC AND DQDOTA ARE SET = DB + QC + SUM FOR I = 0 TO N-1 OF
C DX(LX+I*INCX) * DY(LY+I*INCY), WHERE QC IS AN EXTENDED
C PRECISION RESULT PREVIOUSLY COMPUTED BY DQDOTI OR DQDOTA
C AND LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N, AND LY IS
C DEFINED IN A SIMILAR WAY USING INCY. THE MP PACKAGE BY
C RICHARD P. BRENT IS USED FOR THE EXTENDED PRECISION ARITHMETIC.
C
C FRED T. KROGH, JPL, 1977, JUNE 1
C2
DOUBLE PRECISION DX(1), DY(1), DB
INTEGER QC(10), QX(10), QY(10)
C
C THE COMMON BLOCK FOR THE MP PACKAGE (MODIFIED TO GIVE IT A NAME)
COMMON /MPCOM/ MPB, MPT, MPM, MPLUN, MPMXR, MPR(12)
DATA I1 / 0 /
C
C IF I1 IS 0 THE MP PACKAGE MUST BE INITIALIZED (MPBLAS SETS I1 = 1)
IF (I1 .EQ. 0) CALL MPBLAS(I1)
IF (DB .EQ. 0.D0) GO TO 20
CALL MPCDM(DB, QX)
CALL MPADD(QC, QX, QC)
20 IF (N .EQ. 0) GO TO 40
IX = 1
IY = 1
IF (INCX .LT. 0) IX = (-N + 1) * INCX + 1
IF (INCY .LT. 0) IY = (-N + 1) * INCY + 1
DO 30 I = 1,N
  CALL MPCDM(DX(IX), QX)
  CALL MPCDM(DY(IY), QY)
  CALL MPMUL(QX, QY, QX)
  CALL MPADD(QC, QX, QC)
  IX = IX + INCX
  IY = IY + INCY
30 CONTINUE
40 CALL MPCMD(QC, DQDOTA)
RETURN
END

COMPLEX FUNCTION CDOTC(N,CX,INCX,CY,INCY)
C
C RETURNS THE DOT PRODUCT FOR COMPLEX CX AND CY, USES CONJUGATE(CX)
C
C CDOTC = SUM FOR I = 0 TO N-1 OF CONJ(CX(LX+I*INCX))*CY(LY+I*INCY),
C WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N, AND LY IS
C DEFINED IN A SIMILAR WAY USING INCY.
C
C COMPLEX CX(1),CY(1)
C
C CDOTC = (0.,0.)
IF(N .LE. 0)RETURN
IF(INCX.EQ.INCY.AND.INCX.GT.0) GO TO 20

```

DD18760

DD18780

DD18790

DD18800

DD18810

DD18820

DD18830

DD18840

DD18850

DD18860

DD18870

DD18880

DD18890

DD18900

DD18910

DD18920

DD18930

DD18940

DD18950

DD18960

DD18970

DD18980

DD18760

DD18780

DD18790

DD18800

DD18810

DD18820

DD18830

DD18840

DD18850

DD18860

DD18870

DD18880

DD18890

DD18900

DD18910

DD18920

DD18930

DD18940

DD18950

DD18960

DD18970

DD18980

DD18760

DD18780

DD18790

DD18800

DD18810

DD18820

DD18830

DD18840

DD18850

DD18860

DD18870

DD18880

DD18890

DD18900

DD18910

DD18920

DD18930

DD18940

DD18950

DD18960

DD18970

DD18980

DD18760

DD18780

DD18790

DD18800

DD18810

DD18820

DD18830

DD18840

DD18850

DD18860

DD18870


```

      KX = 1
      KY = 1
      IF(INCX.LT.0) KX = 1+(1-N)*INCX
      IF(INCY.LT.0) KY = 1+(1-N)*INCY
      DO 10 I = 1,N
        CDOTC = CDOTC + CONJG(CX(KX))*CY(KY)
        KX = KX + INCX
        KY = KY + INCY
10    CONTINUE
      RETURN
20  CONTINUE
      NS = N*INCX
      DO 30 I=1,NS,INCX
        CDOTC = CONJG(CX(I))*CY(I) + CDOTC
30  CONTINUE
      RETURN
      END

```

CD19870
CD19880
CD19890
CD19900
CD19910
CD19920
CD19930
CD19940
CD19950
CD19960
CD19970
CD19980
CD19990
CD20000
CD20010
CD20020
CD20030

COMPLEX FUNCTION CDOTU(N,CX,INCX,CY,INCY) CD20040

C
C RETURNS THE DOT PRODUCT FOR COMPLEX CX AND CY, NO CONJUGATION
C CDOTU = SUM FOR I = 0 TO N-1 OF CX(LX+I*INCX) * CY(LY+I*INCY),
C WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N, AND LY IS
C DEFINED IN A SIMILAR WAY USING INCY.
C

COMPLEX CX(1),CY(1) CD20080

C
C CDOTU = (0.,0.)
C IF(N .LE. 0)RETURN CD20090
C IF(INCX.EQ.INCY.AND.INCX.GT.0) GO TO 20 CD20100
C KX = 1 CD20110
C KY = 1 CD20120
C IF(INCX.LT.0) KX = 1+(1-N)*INCX CD20130
C IF(INCY.LT.0) KY = 1+(1-N)*INCY CD20140
C DO 10 I = 1,N CD20150
C CDOTU = CDOTU + CX(KX)*CY(KY) CD20160
C KX = KX + INCX CD20170
C KY = KY + INCY CD20180
10 CONTINUE CD20190
C RETURN CD20200
20 CONTINUE CD20210
C NS = N*INCX CD20220
C DO 30 I=1,NS,INCX CD20230
C CDOTU = CDOTU + CX(I)*CY(I) CD20240
30 CONTINUE CD20250
C RETURN CD20260
C END CD20270
CD20280
CD20290

SUBROUTINE SAXPY(N,SA,SX,INCX,SY,INCY) SA20300

C
C OVERWRITE SINGLE PRECISION SY WITH SINGLE PRECISION SA*SX +SY. SA20310
C FOR I = 0 TO N-1, REPLACE SY(LY+I*INCY) WITH SA*SX(LX+I*INCX) + SA20320
C SY(LY+I*INCY), WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N,
C AND LY IS DEFINED IN A SIMILAR WAY USING INCY.
C

REAL SX(1),SY(1),SA SA20330
IF(N.LE.0.OR.SA.EQ.0.E0) RETURN SA20340
IF(INCX.EQ.INCY) IF(INCX-1) 5,20,60 SA20350
5 CONTINUE SA20360
SA20370
SA20380
SA20390
SA20400
SA20410
SA20420
SA20430
SA20440
SA20450
SA20460
SA20470
SA20480
SA20490
SA20500

C
C CODE FOR NONEQUAL OR NONPOSITIVE INCREMENTS.
C
C IX = 1
C IY = 1
C IF(INCX.LT.0) IX = (-N+1)*INCX + 1
C IF(INCY.LT.0) IY = (-N+1)*INCY + 1
C DO 10 I = 1,N
C SY(IY) = SY(IY) + SA*SX(IX)
C IX = IX + INCX
C IY = IY + INCY
10 CONTINUE
C RETURN

```

C
C          CODE FOR BOTH INCREMENTS EQUAL TO 1
C
C
C          CLEAN-UP LOOP SO REMAINING VECTOR LENGTH IS A MULTIPLE OF 4.
C
20 M = MOD(N,4)
IF( M .EQ. 0 ) GO TO 40
DO 30 I = 1,M
  SY(I) = SY(I) + SA*SX(I)
30 CONTINUE
IF( N .LT. 4 ) RETURN
40 MP1 = M + 1
DO 50 I = MP1,N,4
  SY(I) = SY(I) + SA*SX(I)
  SY(I + 1) = SY(I + 1) + SA*SX(I + 1)
  SY(I + 2) = SY(I + 2) + SA*SX(I + 2)
  SY(I + 3) = SY(I + 3) + SA*SX(I + 3)
50 CONTINUE
RETURN

C
C          CODE FOR EQUAL, POSITIVE, NONUNIT INCREMENTS.
C
60 CONTINUE
NS = N*INCX
DO 70 I=1,NS,INCX
  SY(I) = SA*SX(I) + SY(I)
70 CONTINUE
RETURN
END

SUBROUTINE DAXPY(N,DA,DX,INCX,DY,INCY)
C
C OVERWRITE DOUBLE PRECISION DY WITH DOUBLE PRECISION DA*DX + DY.
C FOR I = 0 TO N-1, REPLACE DY(LY+I*INCY) WITH DA*DX(LX+I*INCX) +
C DY(LY+I*INCY), WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N,
C AND LY IS DEFINED IN A SIMILAR WAY USING INCY.
C
DOUBLE PRECISION DX(1),DY(1),DA
IF(N.LE.0.OR.DA.EQ.0.D0) RETURN
IF(INCX.EQ.INCY) IF(INCX-1) 5,20,60
5 CONTINUE

C
C          CODE FOR NONEQUAL OR NONPOSITIVE INCREMENTS.
C
IX = 1
IY = 1
IF(INCX.LT.0)IX = (-N+1)*INCX + 1
IF(INCY.LT.0)IY = (-N+1)*INCY + 1
DO 10 I = 1,N
  DY(IY) = DY(IY) + DA*DX(IX)
  IX = IX + INCX
  IY = IY + INCY
10 CONTINUE
RETURN

C
C          CODE FOR BOTH INCREMENTS EQUAL TO 1
C
C
C          CLEAN-UP LOOP SO REMAINING VECTOR LENGTH IS A MULTIPLE OF 4.
C
20 M = MOD(N,4)
IF( M .EQ. 0 ) GO TO 40
DO 30 I = 1,M

  DY(I) = DY(I) + DA*DX(I)
30 CONTINUE
IF( N .LT. 4 ) RETURN
40 MP1 = M + 1
DO 50 I = MP1,N,4
  DY(I) = DY(I) + DA*DX(I)
  DY(I + 1) = DY(I + 1) + DA*DX(I + 1)
  DY(I + 2) = DY(I + 2) + DA*DX(I + 2)
  DY(I + 3) = DY(I + 3) + DA*DX(I + 3)
50 CONTINUE

```

```

SA20510
SA20520
SA20530
SA20540
SA20550
SA20560
SA20580
SA20590
SA20600
SA20610
SA20620
SA20630
SA20640
SA20650
SA20660
SA20670
SA20680
SA20690
SA20700
SA20710
SA20720
SA20730
SA20740
SA20750
SA20760
SA20770
SA20780
SA20790
SA20800
DA20810
DA20820
DA20830
DA20840
DA20850
DA20860
DA20870
DA20880
DA20890
DA20900
DA20910
DA20920
DA20930
DA20940
DA20950
DA20960
DA20970
DA20980
DA20990
DA21000
DA21010
DA21020
DA21030
DA21040
DA21050
DA21060
DA21070
DA21090
DA21100
DA21110
DA21120
DA21130
DA21140
DA21150
DA21160
DA21170
DA21180
DA21190
DA21200

```

```

      RETURN
C
C      CODE FOR EQUAL, POSITIVE, NONUNIT INCREMENTS.
C
60 CONTINUE
   NS = N*INCX
      DO 70 I=1,NS,INCX
         DY(I) = DA*DX(I) + DY(I)
70 CONTINUE
   RETURN
   END
DA21210
DA21220
DA21230
DA21240
DA21250
DA21260
DA21270
DA21280
DA21290
DA21300
DA21310

      SUBROUTINE CAXPY(N,CA,CX,INCX,CY,INCY)
C
C      OVERWRITE COMPLEX CY WITH COMPLEX CA*CX + CY.
C      FOR I = 0 TO N-1, REPLACE CY(LY+I*INCY) WITH CA*CX(LX+I*INCX) +
C      CY(LY+I*INCY), WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N,
C      AND LY IS DEFINED IN A SIMILAR WAY USING INCY.
C
      COMPLEX CX(1),CY(1),CA
C
      CANORM = ABS(REAL(CA)) + ABS(AIMAG(CA))
      IF(N.LE.0.OR.CANORM.EQ.0.E0) RETURN
      IF(INCX.EQ.INCY.AND.INCX.GT.0) GO TO 20
      KX = 1
      KY = 1
      IF(INCX.LT.0) KX = 1+(1-N)*INCX
      IF(INCY.LT.0) KY = 1+(1-N)*INCY
      DO 10 I = 1,N
         CY(KY) = CY(KY) + CA*CX(KX)
         KX = KX + INCX
         KY = KY + INCY
10 CONTINUE
      RETURN
20 CONTINUE
   NS = N*INCX
      DO 30 I=1,NS,INCX
         CY(I) = CA*CX(I) + CY(I)
30 CONTINUE
   RETURN
   END
CA21320
CA21360
CA21370
CA21380
CA21390
CA21400
CA21410
CA21420
CA21430
CA21440
CA21450
CA21460
CA21470
CA21480
CA21490
CA21500
CA21510
CA21520
CA21530
CA21540
CA21550
CA21560
CA21570

      SUBROUTINE SROTG(SA,SB,SC,SS)
C
C      DESIGNED BY C.L.LAWSON, JPL, 1977 SEPT 08
C
C      CONSTRUCT THE GIVENS TRANSFORMATION
C
C      ( SC  SS )
C      G = (     ),   SC**2 + SS**2 = 1 ,
C      (-SS  SC )
C
C      WHICH ZEROS THE SECOND ENTRY OF THE 2-VECTOR (SA,SB)**T .
C
C      THE QUANTITY R = (+/-)SQRT(SA**2 + SB**2) OVERWRITES SA IN
C      STORAGE. THE VALUE OF SB IS OVERWRITTEN BY A VALUE Z WHICH
C      ALLOWS SC AND SS TO BE RECOVERED BY THE FOLLOWING ALGORITHM:
C      IF Z=1 SET SC=0. AND SS=1.
C      IF ABS(Z) .LT. 1 SET SC=SQRT(1-Z**2) AND SS=Z
C      IF ABS(Z) .GT. 1 SET SC=1/Z AND SS=SQRT(1-SC**2)
C
C      NORMALLY, THE SUBPROGRAM SROT(N,SX,INCX,SY,INCY,SC,SS) WILL
C      NEXT BE CALLED TO APPLY THE TRANSFORMATION TO A 2 BY N MATRIX.
C
C      -----
C
C      IF (ABS(SA) .LE. ABS(SB)) GO TO 10
C
C      *** HERE ABS(SA) .GT. ABS(SB) ***
C
      U = SA + SA
      V = SB / U

```

```

C
C   NOTE THAT U AND R HAVE THE SIGN OF SA
C
C   R = SQRT(.25 + V**2) * U
C
C   NOTE THAT SC IS POSITIVE
C
C   SC = SA / R
C   SS = V * (SC + SC)
C   SB = SS
C   SA = R
C   RETURN
C
C *** HERE ABS(SA) .LE. ABS(SB) ***
C
C 10 IF (SB .EQ. 0.) GO TO 20
C   U = SB + SB
C   V = SA / U
C
C   NOTE THAT U AND R HAVE THE SIGN OF SB
C   (R IS IMMEDIATELY STORED IN SA)
C
C   SA = SQRT(.25 + V**2) * U
C
C   NOTE THAT SS IS POSITIVE
C
C   SS = SB / SA
C   SC = V * (SS + SS)
C   IF (SC .EQ. 0.) GO TO 15
C   SB = 1. / SC
C   RETURN
C 15 SB = 1.
C   RETURN
C
C *** HERE SA = SB = 0. ***
C
C 20 SC = 1.
C   SS = 0.
C   RETURN
C
C   END

SUBROUTINE DROTG(DA,DB,DC,DS)
C
C   DESIGNED BY C.L.LAWSON, JPL, 1977 SEPT 08
C
C   CONSTRUCT THE GIVENS TRANSFORMATION
C
C       ( DC  DS )
C   G = (      ) ,   DC**2 + DS**2 = 1 ,
C       (-DS  DC )
C
C   WHICH ZEROS THE SECOND ENTRY OF THE 2-VECTOR (DA,DB)**T .
C
C   THE QUANTITY R = (+/-)DSQRT(DA**2 + DB**2) OVERWRITES DA IN
C   STORAGE. THE VALUE OF DB IS OVERWRITTEN BY A VALUE Z WHICH
C   ALLOWS DC AND DS TO BE RECOVERED BY THE FOLLOWING ALGORITHM:
C   IF Z=1 SET DC=0.D0 AND DS=1.D0
C   IF DABS(Z) .LT. 1 SET DC=DSQRT(1-Z**2) AND DS=Z
C   IF DABS(Z) .GT. 1 SET DC=1/Z AND DS=DSQRT(1-DC**2)
C
C   NORMALLY, THE SUBPROGRAM DROT(N,DX,INCX,DY,INCY,DC,DS) WILL
C   NEXT BE CALLED TO APPLY THE TRANSFORMATION TO A 2 BY N MATRIX.
C
C -----
C
C   DOUBLE PRECISION DA, DB, DC, DS, U, V, R
C   IF (DABS(DA) .LE. DABS(DB)) GO TO 10
C
C *** HERE DABS(DA) .GT. DABS(DB) ***
C

```

```

      U = DA + DA
      V = DB / U
C
C   NOTE THAT U AND R HAVE THE SIGN OF DA
C
      R = DSQRT(.25D0 + V**2) * U
C
C   NOTE THAT DC IS POSITIVE
C
      DC = DA / R
      DS = V * (DC + DC)
      DB = DS
      DA = R
      RETURN
C
C *** HERE DABS(DA) .LE. DABS(DB) ***
C
      10 IF (DB .EQ. 0.D0) GO TO 20
         U = DB + DB
         V = DA / U
C
C   NOTE THAT U AND R HAVE THE SIGN OF DB
C   (R IS IMMEDIATELY STORED IN DA)
C
      DA = DSQRT(.25D0 + V**2) * U
C
C   NOTE THAT DS IS POSITIVE
C
      DS = DB / DA
      DC = V * (DS + DS)
      IF (DC .EQ. 0.D0) GO TO 15
      DB = 1.D0 / DC
      RETURN
      15 DB = 1.D0
         RETURN
C
C *** HERE DA = DB = 0.D0 ***
C
      20 DC = 1.D0
         DS = 0.D0
         RETURN
C
      END

      SUBROUTINE SROT(N,SX,INCX,SY,INCY,SC,SS)
C
C   MULTIPLY THE 2 X 2 MATRIX ( SC SS) TIMES THE 2 X N MATRIX (SX**T)
C   (-SS SC) (SY**T)
C   WHERE **T INDICATES TRANSPOSE. THE ELEMENTS OF SX ARE IN
C
C   SX(LX+I*INCX), I = 0 TO N-1, WHERE LX = 1 IF INCX .GE. 0, ELSE
C   LX = (-INCX)*N, AND SIMILARLY FOR SY USING LY AND INCY.
      REAL SX,SY,SC,SS,ZERO,ONE,W,Z
      DIMENSION SX(1),SY(1)
C
      DATA ZERO,ONE/0.E0,1.E0/
      IF(N .LE. 0 .OR. (SS .EQ. ZERO .AND. SC .EQ. ONE)) GO TO 40
      IF(.NOT. (INCX .EQ. INCY .AND. INCX .GT. 0)) GO TO 20
C
      NSTEPS=INCX*N
      DO 10 I=1,NSTEPS,INCX
         W=SX(I)
         Z=SY(I)
         SX(I)=SC*W+SS*Z
         SY(I)=-SS*W+SC*Z
      10 CONTINUE
      GO TO 40
C
      20 CONTINUE
      KX=1
      KY=1
C
      IF(INCX .LT. 0) KX=1-(N-1)*INCX
      IF(INCY .LT. 0) KY=1-(N-1)*INCY

```

SR24210

SR24220

SR24260

SR24270

SR24280

SR24290

SR24300

SR24310

SR24320

SR24330

SR24340

SR24350

SR24360

SR24370

SR24380

SR24390

SR24400

SR24410

SR24440

SR24450

SR24460

SR24470

SR24480

```

C
      DO 30 I=1,N
          W=DX(KX)
          Z=DY(KY)
          SX(KX)=SC*W+SS*Z
          SY(KY)=-SS*W+SC*Z
          KX=KX+INCX
          KY=KY+INCY
      30 CONTINUE
40 CONTINUE
C
      RETURN
      END
C
      SUBROUTINE DROT(N,DX,INCX,DY,INCY,DC,DS)
C
C      MULTIPLY THE 2 X 2 MATRIX ( DC DS) TIMES THE 2 X N MATRIX (DX**T)
C      (-DS DC) (DY**T)
C      WHERE **T INDICATES TRANSPOSE. THE ELEMENTS OF DX ARE IN
C      DX(LX+I*INCX), I = 0 TO N-1, WHERE LX = 1 IF INCX .GE. 0, ELSE
C      LX = (-INCX)*N, AND SIMILARLY FOR DY USING LY AND INCY.
C      DOUBLE PRECISION DX,DY,DC,DS,ZERO,ONE,W,Z
      DIMENSION DX(1),DY(1)
C
      DATA ZERO,ONE/0.D0,1.D0/
      IF(N .LE. 0 .OR. (DS .EQ. ZERO .AND. DC .EQ. ONE)) GO TO 40
      IF(.NOT. (INCX .EQ. INCY .AND. INCX .GT. 0)) GO TO 20
C
      NSTEPS=INCX*N
      DO 10 I=1,NSTEPS,INCX
          W=DX(I)
          Z=DY(I)
          DX(I)=DC*W+DS*Z
          DY(I)=-DS*W+DC*Z
      10 CONTINUE
      GO TO 40
C
      20 CONTINUE
      KX=1
      KY=1
C
      IF(INCX .LT. 0) KX=1-(N-1)*INCX
      IF(INCY .LT. 0) KY=1-(N-1)*INCY
C
      DO 30 I=1,N
          W=DX(KX)
          Z=DY(KY)
          DX(KX)=DC*W+DS*Z
          DY(KY)=-DS*W+DC*Z
          KX=KX+INCX
          KY=KY+INCY
      30 CONTINUE
40 CONTINUE
C
      RETURN
      END
C
      SUBROUTINE SROTMG (SD1,SD2, SX1,SY1,SPARAM)
C
C      CONSTRUCT THE MODIFIED GIVENS TRANSFORMATION MATRIX H WHICH ZEROS
C      THE SECOND COMPONENT OF THE 2-VECTOR (SQRT(SD1)*SX1,SQRT(SD2)*
C      SY2)**T.
C      WITH SPARAM(1)=SFLAG, H HAS ONE OF THE FOLLOWING FORMS..
C
C      SFLAG=-1.E0      SFLAG=0.E0      SFLAG=1.E0      SFLAG=-2.E0
C
C      (SH11 SH12) (1.E0 SH12) (SH11 1.E0) (1.E0 0.E0)
C      H=( ) ( ) ( ) ( )
C      (SH21 SH22), (SH21 1.E0), (-1.E0 SH22), (0.E0 1.E0).
C      LOCATIONS 2-4 OF SPARAM CONTAIN SH11,SH21,SH12, AND SH22
C      RESPECTIVELY. (VALUES OF 1.E0, -1.E0, OR 0.E0 IMPLIED BY THE

```

C	VALUE OF SPARAM(1) ARE NOT STORED IN SPARAM.)	SRO00150
C		SRO00160
C	THE VALUES OF GAMSQ AND RGAMSQ SET IN THE DATA STATEMENT MAY BE	SRO00170
C	INEXACT. THIS IS OK AS THEY ARE ONLY USED FOR TESTING THE SIZE	SRO00180
C	OF SD1 AND SD2. ALL ACTUAL SCALING OF DATA IS DONE USING GAM.	SRO00190
C		SRO00200
C	DIMENSION SPARAM(5)	SRO00210
C		SRO00220
	DATA ZERO,ONE,TWO /0.E0,1.E0,2.E0/	SRO00230
	DATA GAM,GAMSQ,RGAMSQ/4096.E0,1.67772E7,5.96046E-8/	SRO00240
	IF(.NOT. SD1 .LT. ZERO) GO TO 10	SRO00250
C	GO ZERO-H-D-AND-SX1..	SRO00260
	GO TO 60	SRO00270
10	CONTINUE	SRO00280
C	CASE-SD1-NONNEGATIVE	SRO00290
	SP2=SD2*SY1	SRO00300
	IF(.NOT. SP2 .EQ. ZERO) GO TO 20	SRO00310
	SFLAG=-TWO	SRO00320
	GO TO 260	SRO00330
C	REGULAR-CASE..	SRO00340
20	CONTINUE	SRO00350
	SP1=SD1*SX1	SRO00360
	SQ2=SP2*SY1	SRO00370
	SQ1=SP1*SX1	SRO00380
C		SRO00390
	IF(.NOT. ABS(SQ1) .GT. ABS(SQ2)) GO TO 40	SRO00400
	SH21=-SY1/SX1	SRO00410
	SH12=SP2/SP1	SRO00420
C		SRO00430
	SU=ONE-SH12*SH21	SRO00440
C		SRO00450
	IF(.NOT. SU .LE. ZERO) GO TO 30	SRO00460
C	GO ZERO-H-D-AND-SX1..	SRO00470
	GO TO 60	SRO00480
30	CONTINUE	SRO00490
	SFLAG=ZERO	SRO00500
	SD1=SD1/SU	SRO00510
	SD2=SD2/SU	SRO00520
	SX1=SX1*SU	SRO00530
C	GO SCALE-CHECK..	SRO00540
	GO TO 100	SRO00550
40	CONTINUE	SRO00560
	IF(.NOT. SQ2 .LT. ZERO) GO TO 50	SRO00570
C	GO ZERO-H-D-AND-SX1..	SRO00580
	GO TO 60	SRO00590
50	CONTINUE	SRO00600
	SFLAG=ONE	SRO00610
	SH11=SP1/SP2	SRO00620
	SH22=SX1/SY1	SRO00630
	SU=ONE+SH11*SH22	SRO00640
	STEMP=SD2/SU	SRO00650
	SD2=SD1/SU	SRO00660
	SD1=STEMP	SRO00670
	SX1=SY1*SU	SRO00680
C	GO SCALE-CHECK	SRO00690
	GO TO 100	SRO00700
C	PROCEDURE..ZERO-H-D-AND-SX1..	SRO00710
60	CONTINUE	SRO00720
	SFLAG=-ONE	SRO00730
	SH11=ZERO	SRO00740
	SH12=ZERO	SRO00750
	SH21=ZERO	SRO00760
	SH22=ZERO	SRO00770
C		SRO00780
	SD1=ZERO	SRO00790
	SD2=ZERO	SRO00800
	SX1=ZERO	SRO00810
C	RETURN..	SRO00820
	GO TO 220	SRO00830
C	PROCEDURE..FIX-H..	SRO00840
70	CONTINUE	SRO00850
	IF(.NOT. SFLAG .GE. ZERO) GO TO 90	SRO00860
C		SRO00870
	IF(.NOT. SFLAG .EQ. ZERO) GO TO 80	SRO00880
	SH11=ONE	SRO00890
	SH22=ONE	SRO00900

	SFLAG=-ONE	SRO00910
	GO TO 90	SRO00920
80	CONTINUE	SRO00930
	SH21=-ONE	SRO00940
	SH12=ONE	SRO00950
	SFLAG=-ONE	SRO00960
90	CONTINUE	SRO00970
	GO TO IGO, (120, 150, 180, 210)	SRO00980
C	PROCEDURE..SCALE-CHECK	SRO00990
100	CONTINUE	SRO01000
110	CONTINUE	SRO01010
	IF(.NOT. SD1 .LE. RGAMSQ) GO TO 130	SRO01020
	IF(SD1 .EQ. ZERO) GO TO 160	SRO01030
	ASSIGN 120 TO IGO	SRO01040
C	FIX-H..	SRO01050
	GO TO 70	SRO01060
120	CONTINUE	SRO01070
	SD1=SD1*GAM**2	SRO01080
	SX1=SD1/GAM	SRO01090
	SH11=SH11/GAM	SRO01100
	SH12=SH12/GAM	SRO01110
	GO TO 110	SRO01120
130	CONTINUE	SRO01130
140	CONTINUE	SRO01140
	IF(.NOT. SD1 .GE. GAMSQ) GO TO 160	SRO01150
	ASSIGN 150 TO IGO	SRO01160
C	FIX-H..	SRO01170
	GO TO 70	SRO01180
150	CONTINUE	SRO01190
	SD1=SD1/GAM**2	SRO01200
	SX1=SD1*GAM	SRO01210
	SH11=SH11*GAM	SRO01220
	SH12=SH12*GAM	SRO01230
	GO TO 140	SRO01240
160	CONTINUE	SRO01250
170	CONTINUE	SRO01260
	IF(.NOT. ABS(SD2) .LE. RGAMSQ) GO TO 190	SRO01270
	IF(SD2 .EQ. ZERO) GO TO 220	SRO01280
	ASSIGN 180 TO IGO	SRO01290
C	FIX-H..	SRO01300
	GO TO 70	SRO01310
180	CONTINUE	SRO01320
	SD2=SD2*GAM**2	SRO01330
	SH21=SH21/GAM	SRO01340
	SH22=SH22/GAM	SRO01350
	GO TO 170	SRO01360
190	CONTINUE	SRO01370
200	CONTINUE	SRO01380
	IF(.NOT. ABS(SD2) .GE. GAMSQ) GO TO 220	SRO01390
	ASSIGN 210 TO IGO	SRO01400
C	FIX-H..	SRO01410
	GO TO 70	SRO01420
210	CONTINUE	SRO01430
	SD2=SD2/GAM**2	SRO01440
	SH21=SH21*GAM	SRO01450
	SH22=SH22*GAM	SRO01460
	GO TO 200	SRO01470
220	CONTINUE	SRO01480
	IF(SFLAG) 250, 230, 240	SRO01490
230	CONTINUE	SRO01500
	SPARAM(3)=SH21	SRO01510
	SPARAM(4)=SH12	SRO01520
	GO TO 260	SRO01530
240	CONTINUE	SRO01540
	SPARAM(2)=SH11	SRO01550
	SPARAM(5)=SH22	SRO01560
	GO TO 260	SRO01570
250	CONTINUE	SRO01580
	SPARAM(2)=SH11	SRO01590
	SPARAM(3)=SH21	SRO01600
	SPARAM(4)=SH12	SRO01610
	SPARAM(5)=SH22	SRO01620
260	CONTINUE	SRO01630
	SPARAM(1)=SFLAG	SRO01640
	RETURN	SRO01650
	END	SRO01660


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SUBROUTINE DROTMG (DD1,DD2,DX1,DY1,DPARAM)                                DRO00010
C                                                                 DRO00020
C   CONSTRUCT THE MODIFIED GIVENS TRANSFORMATION MATRIX H WHICH ZEROS DRO00030
C   THE SECOND COMPONENT OF THE 2-VECTOR (DSQRT(DD1)*DX1,DSQRT(DD2))* DRO00040
C   DY2)**T. DRO00050
C   WITH DPARAM(1)=DFLAG, H HAS ONE OF THE FOLLOWING FORMS.. DRO00060
C                                                                 DRO00070
C   DFLAG=-1.D0      DFLAG=0.D0      DFLAG=1.D0      DFLAG=-2.D0      DRO00080
C                                                                 DRO00090
C   (DH11 DH12)      (1.D0 DH12)      (DH11 1.D0)      (1.D0 0.D0)      DRO00100
C   H=(              ) (              ) (              ) (              ) DRO00110
C   (DH21 DH22),      (DH21 1.D0),      (-1.D0 DH22),      (0.D0 1.D0). DRO00120
C   LOCATIONS 2-4 OF DPARAM CONTAIN DH11, DH21, DH12, AND DH22 DRO00130
C   RESPECTIVELY. (VALUES OF 1.D0, -1.D0, OR 0.D0 IMPLIED BY THE DRO00140
C   VALUE OF DPARAM(1) ARE NOT STORED IN DPARAM.) DRO00150
C                                                                 DRO00160
C   THE VALUES OF GAMSQ AND RGAMSQ SET IN THE DATA STATEMENT MAY BE DRO00170
C   INEXACT. THIS IS OK AS THEY ARE ONLY USED FOR TESTING THE SIZE DRO00180
C   OF DD1 AND DD2. ALL ACTUAL SCALING OF DATA IS DONE USING GAM. DRO00190
C                                                                 DRO00200
C   DOUBLE PRECISION GAM,ONE,RGAMSQ,DD2,DH11,DH21,DPARAM,DP2, DRO00210
C   1 DQ2,DU,DY1,ZERO,GAMSQ,DD1,DFLAG,DH12,DH22,DP1,DQ1, DRO00220
C   2 DTEMP,DX1,TWO DRO00230
C   DIMENSION DPARAM(5) DRO00240
C                                                                 DRO00250
C   DATA ZERO,ONE,TWO /0.D0,1.D0,2.D0/ DRO00260
C   DATA GAM,GAMSQ,RGAMSQ/4096.D0,16777216.D0,5.9604645D-8/ DRO00270
C   IF(.NOT. DD1 .LT. ZERO) GO TO 10 DRO00280
C   GO ZERO-H-D-AND-DX1.. DRO00290
C   GO TO 60 DRO00300
10 CONTINUE DRO00310
C   CASE-DD1-NONNEGATIVE DRO00320
C   DP2=DD2*DY1 DRO00330
C   IF(.NOT. DP2 .EQ. ZERO) GO TO 20 DRO00340
C   DFLAG=-TWO DRO00350
C   GO TO 260 DRO00360
C   REGULAR-CASE.. DRO00370
20 CONTINUE DRO00380
C   DP1=DD1*DX1 DRO00390
C   DQ2=DP2*DY1 DRO00400
C   DQ1=DP1*DX1 DRO00410
C                                                                 DRO00420
C   IF(.NOT. DABS(DQ1) .GT. DABS(DQ2)) GO TO 40 DRO00430
C   DH21=-DY1/DX1 DRO00440
C   DH12=DP2/DP1 DRO00450
C                                                                 DRO00460
C   DU=ONE-DH12*DH21 DRO00470
C                                                                 DRO00480
C   IF(.NOT. DU .LE. ZERO) GO TO 30 DRO00490
C   GO ZERO-H-D-AND-DX1.. DRO00500
C   GO TO 60 DRO00510
30 CONTINUE DRO00520
C   DFLAG=ZERO DRO00530
C   DD1=DD1/DU DRO00540
C   DD2=DD2/DU DRO00550
C   DX1=DX1*DU DRO00560
C   GO SCALE-CHECK.. DRO00570
C   GO TO 100 DRO00580
40 CONTINUE DRO00590
C   IF(.NOT. DQ2 .LT. ZERO) GO TO 50 DRO00600
C   GO ZERO-H-D-AND-DX1.. DRO00610
C   GO TO 60 DRO00620
50 CONTINUE DRO00630
C   DFLAG=ONE DRO00640
C   DH11=DP1/DP2 DRO00650
C   DH22=DX1/DY1 DRO00660
C   DU=ONE+DH11*DH22 DRO00670
C   DTEMP=DD2/DU DRO00680
C   DD2=DD1/DU DRO00690
C   DD1=DTEMP DRO00700
C   DX1=DY1*DU DRO00710
C   GO SCALE-CHECK DRO00720
C   GO TO 100 DRO00730
C   PROCEDURE..ZERO-H-D-AND-DX1.. DRO00740
60 CONTINUE DRO00750
C   DFLAG=-ONE DRO00760

```

	DH11=ZERO	DRO00770
	DH12=ZERO	DRO00780
	DH21=ZERO	DRO00790
	DH22=ZERO	DRO00800
C		DRO00810
	DD1=ZERO	DRO00820
	DD2=ZERO	DRO00830
	DX1=ZERO	DRO00840
C	RETURN..	DRO00850
	GO TO 220	DRO00860
C	PROCEDURE..FIX-H..	DRO00870
70	CONTINUE	DRO00880
	IF(.NOT. DFLAG .GE. ZERO) GO TO 90	DRO00890
C		DRO00900
	IF(.NOT. DFLAG .EQ. ZERO) GO TO 80	DRO00910
	DH11=ONE	DRO00920
	DH22=ONE	DRO00930
	DFLAG=-ONE	DRO00940
	GO TO 90	DRO00950
80	CONTINUE	DRO00960
	DH21=-ONE	DRO00970
	DH12=ONE	DRO00980
	DFLAG=-ONE	DRO00990
90	CONTINUE	DRO01000
	GO TO IGO, (120,150,180,210)	DRO01010
C	PROCEDURE..SCALE-CHECK	DRO01020
100	CONTINUE	DRO01030
110	CONTINUE	DRO01040
	IF(.NOT. DD1 .LE. RGAMSQ) GO TO 130	DRO01050
	IF(DD1 .EQ. ZERO) GO TO 160	DRO01060
	ASSIGN 120 TO IGO	DRO01070
C	FIX-H..	DRO01080
	GO TO 70	DRO01090
120	CONTINUE	DRO01100
	DD1=DD1*GAM**2	DRO01110
	DX1=DX1/GAM	DRO01120
	DH11=DH11/GAM	DRO01130
	DH12=DH12/GAM	DRO01140
	GO TO 110	DRO01150
130	CONTINUE	DRO01160
140	CONTINUE	DRO01170
	IF(.NOT. DD1 .GE. GAMSQ) GO TO 160	DRO01180
	ASSIGN 150 TO IGO	DRO01190
C	FIX-H..	DRO01200
	GO TO 70	DRO01210
150	CONTINUE	DRO01220
	DD1=DD1/GAM**2	DRO01230
	DX1=DX1*GAM	DRO01240
	DH11=DH11*GAM	DRO01250
	DH12=DH12*GAM	DRO01260
	GO TO 140	DRO01270
160	CONTINUE	DRO01280
170	CONTINUE	DRO01290
	IF(.NOT. DABS(DD2) .LE. RGAMSQ) GO TO 190	DRO01300
	IF(DD2 .EQ. ZERO) GO TO 220	DRO01310
	ASSIGN 180 TO IGO	DRO01320
C	FIX-H..	DRO01330
	GO TO 70	DRO01340
180	CONTINUE	DRO01350
	DD2=DD2*GAM**2	DRO01360
	DH21=DH21/GAM	DRO01370
	DH22=DH22/GAM	DRO01380
	GO TO 170	DRO01390
190	CONTINUE	DRO01400
200	CONTINUE	DRO01410
	IF(.NOT. DABS(DD2) .GE. GAMSQ) GO TO 220	DRO01420
	ASSIGN 210 TO IGO	DRO01430
C	FIX-H..	DRO01440
	GO TO 70	DRO01450
210	CONTINUE	DRO01460
	DD2=DD2/GAM**2	DRO01470
	DH21=DH21*GAM	DRO01480
	DH22=DH22*GAM	DRO01490
	GO TO 200	DRO01500
220	CONTINUE	DRO01510
	IF(DFLAG) 250, 230, 240	DRO01520

230	CONTINUE	DRO01530
	DPARAM(3)=DH21	DRO01540
	DPARAM(4)=DH12	DRO01550
	GO TO 260	DRO01560
240	CONTINUE	DRO01570
	DPARAM(2)=DH11	DRO01580
	DPARAM(5)=DH22	DRO01590
	GO TO 260	DRO01600
250	CONTINUE	DRO01610
	DPARAM(2)=DH11	DRO01620
	DPARAM(3)=DH21	DRO01630
	DPARAM(4)=DH12	DRO01640
	DPARAM(5)=DH22	DRO01650
260	CONTINUE	DRO01660
	DPARAM(1)=DFLAG	DRO01670
	RETURN	DRO01680
	END	DRO01690
	SUBROUTINE SROTM (N,SX,INCX,SY,INCY,SPARAM)	SR28430
C		SR28440
C	APPLY THE MODIFIED GIVENS TRANSFORMATION, H, TO THE 2 BY N MATRIX	SR28450
C		SR28460
C	(SX**T) , WHERE **T INDICATES TRANSPOSE. THE ELEMENTS OF SX ARE IN	
C	(DX**T)	
C		
C	SX(LX+I*INCX), I = 0 TO N-1, WHERE LX = 1 IF INCX .GE. 0, ELSE	
C	LX = (-INCX)*N, AND SIMILARLY FOR SY USING LY AND INCY.	
C	WITH SPARAM(1)=SFLAG, H HAS ONE OF THE FOLLOWING FORMS..	SR28510
C		SR28520
C	SFLAG=-1.E0 SFLAG=0.E0 SFLAG=1.E0 SFLAG=-2.E0	SR28530
C		SR28540
C	(SH11 SH12) (1.E0 SH12) (SH11 1.E0) (1.E0 0.E0)	SR28550
C	H=() () () ()	SR28560
C	(SH21 SH22), (SH21 1.E0), (-1.E0 SH22), (0.E0 1.E0).	SR28570
C	SEE SROTMG FOR A DESCRIPTION OF DATA STORAGE IN SPARAM.	
C		SR28580
	DIMENSION SX(1),SY(1),SPARAM(5)	SR28590
	DATA ZERO,TWO/0.E0,2.E0/	SR28600
C		SR28610
	SFLAG=SPARAM(1)	SR28620
	IF(N .LE. 0 .OR. (SFLAG+TWO.EQ.ZERO)) GO TO 140	SR28630
	IF(.NOT.(INCX.EQ.INCY.AND. INCX .GT.0)) GO TO 70	SR28640
C		SR28650
	NSTEPS=N*INCX	SR28660
	IF(SFLAG) 50,10,30	SR28670
10	CONTINUE	SR28680
	SH12=SPARAM(4)	SR28690
	SH21=SPARAM(3)	SR28700
	DO 20 I=1,NSTEPS,INCX	SR28710
	W=SX(I)	SR28720
	Z=SY(I)	SR28730
	SX(I)=W+Z*SH12	SR28740
	SY(I)=W*SH21+Z	SR28750
20	CONTINUE	SR28760
	GO TO 140	SR28770
30	CONTINUE	SR28780
	SH11=SPARAM(2)	SR28790
	SH22=SPARAM(5)	SR28800
	DO 40 I=1,NSTEPS,INCX	SR28810
	W=SX(I)	SR28820
	Z=SY(I)	SR28830
	SX(I)=W*SH11+Z	SR28840
	SY(I)=-W+SH22*Z	SR28850
40	CONTINUE	SR28860
	GO TO 140	SR28870
50	CONTINUE	SR28880
	SH11=SPARAM(2)	SR28890
	SH12=SPARAM(4)	SR28900
	SH21=SPARAM(3)	SR28910
	SH22=SPARAM(5)	SR28920
	DO 60 I=1,NSTEPS,INCX	SR28930
	W=SX(I)	SR28940
	Z=SY(I)	SR28950
	SX(I)=W*SH11+Z*SH12	SR28960

	SY(I)=W*SH21+Z*SH22	SR28970
60	CONTINUE	SR28980
	GO TO 140	SR28990
70	CONTINUE	SR29000
	KX=1	SR29010
	KY=1	SR29020
	IF(INCX .LT. 0) KX=1+(1-N)*INCX	SR29030
	IF(INCY .LT. 0) KY=1+(1-N)*INCY	SR29040
C		SR29050
	IF(SFLAG)120,80,100	SR29060
80	CONTINUE	SR29070
	SH12=SPARAM(4)	SR29080
	SH21=SPARAM(3)	SR29090
	DO 90 I=1,N	SR29100
	W=SX(KX)	SR29110
	Z=SY(KY)	SR29120
	SX(KX)=W+Z*SH12	SR29130
	SY(KY)=W*SH21+Z	SR29140
	KX=KX+INCX	SR29150
	KY=KY+INCY	SR29160
90	CONTINUE	SR29170
	GO TO 140	SR29180
100	CONTINUE	SR29190
	SH11=SPARAM(2)	SR29200
	SH22=SPARAM(5)	SR29210
	DO 110 I=1,N	SR29220
	W=SX(KX)	SR29230
	Z=SY(KY)	SR29240
	SX(KX)=W*SH11+Z	SR29250
	SY(KY)=-W+SH22*Z	SR29260
	KX=KX+INCX	SR29270
	KY=KY+INCY	SR29280
110	CONTINUE	SR29290
	GO TO 140	SR29300
120	CONTINUE	SR29310
	SH11=SPARAM(2)	SR29320
	SH12=SPARAM(4)	SR29330
	SH21=SPARAM(3)	SR29340
	SH22=SPARAM(5)	SR29350
	DO 130 I=1,N	SR29360
	W=SX(KX)	SR29370
	Z=SY(KY)	SR29380
	SX(KX)=W*SH11+Z*SH12	SR29390
	SY(KY)=W*SH21+Z*SH22	SR29400
	KX=KX+INCX	SR29410
	KY=KY+INCY	SR29420
130	CONTINUE	SR29430
140	CONTINUE	SR29440
	RETURN	SR29450
	END	SR29460
	SUBROUTINE SCOPY(N,SX,INCX,SY,INCY)	SC30530
C		SC30540
C	COPY SINGLE PRECISION SX TO SINGLE PRECISION SY.	SC30550
C	FOR I = 0 TO N-1, COPY SX(LX+I*INCX) TO SY(LY+I*INCY),	
C	WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N, AND LY IS	
C	DEFINED IN A SIMILAR WAY USING INCY.	
C		SC30560
	REAL SX(1),SY(1)	SC30570
	IF(N.LE.0)RETURN	SC30580
	IF(INCX.EQ.INCY) IF(INCX-1) 5,20,60	SC30590
5	CONTINUE	SC30600
		SC30610
C		SC30620
C	CODE FOR UNEQUAL OR NONPOSITIVE INCREMENTS.	SC30630
C		SC30640
	IX = 1	
	IY = 1	SC30650
	IF(INCX.LT.0)IX = (-N+1)*INCX + 1	SC30660
	IF(INCY.LT.0)IY = (-N+1)*INCY + 1	SC30670
	DO 10 I = 1,N	SC30680
	SY(IY) = SX(IX)	SC30690
	IX = IX + INCX	SC30700

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        IY = IY + INCY
10 CONTINUE
    RETURN
C
C        CODE FOR BOTH INCREMENTS EQUAL TO 1
C
C
C        CLEAN-UP LOOP SO REMAINING VECTOR LENGTH IS A MULTIPLE OF 7.
C
20 M = MOD(N,7)
    IF( M .EQ. 0 ) GO TO 40
    DO 30 I = 1,M
        SY(I) = SX(I)
30 CONTINUE
    IF( N .LT. 7 ) RETURN
40 MP1 = M + 1
    DO 50 I = MP1,N,7
        SY(I) = SX(I)
        SY(I + 1) = SX(I + 1)
        SY(I + 2) = SX(I + 2)
        SY(I + 3) = SX(I + 3)
        SY(I + 4) = SX(I + 4)
        SY(I + 5) = SX(I + 5)
        SY(I + 6) = SX(I + 6)
50 CONTINUE
    RETURN
C
C        CODE FOR EQUAL, POSITIVE, NONUNIT INCREMENTS.
C
60 CONTINUE
    NS = N*INCX
    DO 70 I=1,NS,INCX
        SY(I) = SX(I)
70 CONTINUE
    RETURN
    END

SUBROUTINE DCOPY(N,DX,INCX,DY,INCY)
C
C    COPY DOUBLE PRECISION DX TO DOUBLE PRECISION DY.
C    FOR I = 0 TO N-1, COPY DX(LX+I*INCX) TO DY(LY+I*INCY),
C    WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N, AND LY IS
C    DEFINED IN A SIMILAR WAY USING INCY.
C
C        DOUBLE PRECISION DX(1),DY(1)
C
C        IF(N.LE.0)RETURN
C        IF(INCX.EQ.INCY) IF(INCX-1) 5,20,60
5 CONTINUE
C
C        CODE FOR UNEQUAL OR NONPOSITIVE INCREMENTS.
C
C
C        IX = 1
C        IY = 1
C        IF(INCX.LT.0)IX = (-N+1)*INCX + 1
C        IF(INCY.LT.0)IY = (-N+1)*INCY + 1
C        DO 10 I = 1,N
C            DY(IY) = DX(IX)
C            IX = IX + INCX
C            IY = IY + INCY
10 CONTINUE
    RETURN
C
C        CODE FOR BOTH INCREMENTS EQUAL TO 1
C
C
C        CLEAN-UP LOOP SO REMAINING VECTOR LENGTH IS A MULTIPLE OF 7.
C
20 M = MOD(N,7)
    IF( M .EQ. 0 ) GO TO 40
    DO 30 I = 1,M
        DY(I) = DX(I)
30 CONTINUE
    IF( N .LT. 7 ) RETURN

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```

SC30710
SC30720
SC30730
SC30740
SC30750
SC30760
SC30770
SC30780
SC30790
SC30810
SC30820
SC30830
SC30840
SC30850
SC30860
SC30870
SC30880
SC30890
SC30900
SC30910
SC30920
SC30930
SC30940
SC30950
SC30960
SC30970
SC30980
SC30990
SC31000
SC31010
SC31020
SC31030
SC31040
SC31050
SC31060
DC31070
DC31080
DC31090
DC31100
DC31110
DC31120
DC31130
DC31140
DC31150
DC31160
DC31170
DC31180
DC31190
DC31200
DC31210
DC31220
DC31230
DC31240
DC31250
DC31260
DC31270
DC31280
DC31290
DC31300
DC31310
DC31320
DC31330
DC31350
DC31360
DC31370
DC31380
DC31390

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40 MP1 = M + 1
DO 50 I = MP1,N,7
  DY(I) = DX(I)
  DY(I + 1) = DX(I + 1)
  DY(I + 2) = DX(I + 2)
  DY(I + 3) = DX(I + 3)
  DY(I + 4) = DX(I + 4)
  DY(I + 5) = DX(I + 5)
  DY(I + 6) = DX(I + 6)
50 CONTINUE
RETURN
C
C      CODE FOR EQUAL, POSITIVE, NONUNIT INCREMENTS.
C
60 CONTINUE
NS=N*INCX
DO 70 I=1,NS,INCX
  DY(I) = DX(I)
70 CONTINUE
RETURN
END
DC31400
DC31410
DC31420
DC31430
DC31440
DC31450
DC31460
DC31470
DC31480
DC31490
DC31500
DC31510
DC31520
DC31530
DC31540
DC31550
DC31560
DC31570
DC31580
DC31590
DC31600

SUBROUTINE CCOPY(N,CX,INCX,CY,INCY)
C
C      COPY COMPLEX CX TO COMPLEX CY.
C      FOR I = 0 TO N-1, COPY CX(LX+I*INCX) TO CY(LY+I*INCY),
C      WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N, AND LY IS
C      DEFINED IN A SIMILAR WAY USING INCY.
C
C      COMPLEX CX(1),CY(1)
C
C      IF(N .LE. 0)RETURN
C      IF(INCX.EQ.INCY.AND.INCX.GT.0) GO TO 20
C      KX = 1
C      KY = 1
C      IF(INCX.LT.0) KX = 1+(1-N)*INCX
C      IF(INCY.LT.0) KY = 1+(1-N)*INCY
C      DO 10 I = 1,N
C        CY(KY) = CX(KX)
C        KX = KX + INCX
C        KY = KY + INCY
10 CONTINUE
RETURN
20 CONTINUE
NS = N*INCX
DO 30 I=1,NS,INCX
  CY(I) = CX(I)
30 CONTINUE
RETURN
END
CC31610
CC31650
CC31660
CC31670
CC31680
CC31690
CC31700
CC31710
CC31720
CC31730
CC31740
CC31750
CC31760
CC31770
CC31780
CC31790
CC31800
CC31810
CC31820
CC31830
CC31840
CC31850

SUBROUTINE SSWAP (N,SX,INCX,SY,INCY)
C
C      INTERCHANGE SINGLE PRECISION SX AND SINGLE PRECISION SY.
C      FOR I = 0 TO N-1, INTERCHANGE SX(LX+I*INCX) AND SY(LY+I*INCY),
C      WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N, AND LY IS
C      DEFINED IN A SIMILAR WAY USING INCY.
C
C      REAL SX(1),SY(1),STEMP1,STEMP2,STEMP3
C      IF(N.LE.0)RETURN
C      IF(INCX.EQ.INCY) IF(INCX-1) 5,20,60
5 CONTINUE
C
C      CODE FOR UNEQUAL OR NONPOSITIVE INCREMENTS.
C
C      IX = 1
C      IY = 1
C      IF(INCX.LT.0)IX = (-N+1)*INCX + 1
C      IF(INCY.LT.0)IY = (-N+1)*INCY + 1
C      DO 10 I = 1,N
C        STEMP1 = SX(IX)
SS31860
SS31870
SS31880
SS31890
SS31900
SS31910
SS31920
SS31930
SS31940
SS31950
SS31960
SS31970
SS31980
SS31990
SS32000
SS32010
SS32020

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```

          SX(IX) = SY(IY)
          SY(IY) = STEMP1
          IX = IX + INCX
SS32030
SS32040
SS32050

          IY = IY + INCY
10 CONTINUE
RETURN
C
C      CODE FOR BOTH INCREMENTS EQUAL TO 1
C
C
C      CLEAN-UP LOOP SO REMAINING VECTOR LENGTH IS A MULTIPLE OF 3.
C
20 M = MOD(N,3)
   IF( M .EQ. 0 ) GO TO 40
   DO 30 I = 1,M
       STEMP1 = SX(I)
       SX(I) = SY(I)
       SY(I) = STEMP1
30 CONTINUE
   IF( N .LT. 3 ) RETURN
40 MP1 = M + 1
   DO 50 I = MP1,N,3
       STEMP1 = SX(I)
       STEMP2 = SX(I+1)
       STEMP3 = SX(I+2)
       SX(I) = SY(I)
       SX(I+1) = SY(I+1)
       SX(I+2) = SY(I+2)
       SY(I) = STEMP1
       SY(I+1) = STEMP2
       SY(I+2) = STEMP3
50 CONTINUE
RETURN
60 CONTINUE
C
C      CODE FOR EQUAL, POSITIVE, NONUNIT INCREMENTS.
C
      NS = N*INCX
      DO 70 I=1,NS,INCX
          STEMP1 = SX(I)
          SX(I) = SY(I)
          SY(I) = STEMP1
70 CONTINUE
RETURN
END
SS32160
SS32170
SS32180
SS32190
SS32200
SS32210
SS32220
SS32230
SS32240
SS32250
SS32260
SS32270
SS32280
SS32290
SS32300
SS32310
SS32320
SS32330
SS32340
SS32350
SS32360
SS32370
SS32380
SS32390
SS32400
SS32410
SS32420
SS32430
SS32440
SS32450
SS32460
SS32470

SUBROUTINE DSWAP(N,DX,INCX,DY,INCY)
C
C      INTERCHANGE DOUBLE PRECISION DX AND DOUBLE PRECISION DY.
C      FOR I = 0 TO N-1, INTERCHANGE DX(LX+I*INCX) AND DY(LY+I*INCY),
C      WHERE LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N, AND LY IS
C      DEFINED IN A SIMILAR WAY USING INCY.
C
      DOUBLE PRECISION DX(1),DY(1),DTEMP1,DTEMP2,DTEMP3
DS32480
DS32490
DS32500

      IF(N.LE.0)RETURN
      IF(INCX.EQ.INCY) IF(INCX-1) 5,20,60
5 CONTINUE
DS32510
DS32520

      IF(N.LE.0)RETURN
      IF(INCX.EQ.INCY) IF(INCX-1) 5,20,60
DS32530
DS32540
DS32550
5 CONTINUE
DS32560
C
C      CODE FOR UNEQUAL OR NONPOSITIVE INCREMENTS.
C
      IX = 1
      IY = 1
      IF(INCX.LT.0) IX = (-N+1)*INCX + 1
      IF(INCY.LT.0) IY = (-N+1)*INCY + 1
      DO 10 I = 1,N
          DTEMP1 = DX(IX)
          DX(IX) = DY(IY)
          DY(IY) = DTEMP1
          IX = IX + INCX
          IY = IY + INCY
10 CONTINUE
RETURN
C
DS32570
DS32580
DS32590
DS32600
DS32610
DS32620
DS32630
DS32640
DS32650
DS32660
DS32670
DS32680
DS32690
DS32700
DS32710

```

C	CODE FOR BOTH INCREMENTS EQUAL TO 1	DS32720
C		DS32730
C		DS32740
C	CLEAN-UP LOOP SO REMAINING VECTOR LENGTH IS A MULTIPLE OF 3.	DS32750
C		DS32760
	20 M = MOD(N,3)	
	IF(M .EQ. 0) GO TO 40	DS32780
	DO 30 I = 1,M	DS32790
	DTEMP1 = DX(I)	DS32800
	DX(I) = DY(I)	DS32810
	DY(I) = DTEMP1	DS32820
	30 CONTINUE	DS32830
	IF(N .LT. 3) RETURN	DS32840
	40 MP1 = M + 1	DS32850
	DO 50 I = MP1,N,3	DS32860
	DTEMP1 = DX(I)	DS32870
	DTEMP2 = DX(I+1)	DS32880
	DTEMP3 = DX(I+2)	DS32890
	DX(I) = DY(I)	DS32900
	DX(I+1) = DY(I+1)	DS32910
	DX(I+2) = DY(I+2)	DS32920
	DY(I) = DTEMP1	DS32930
	DY(I+1) = DTEMP2	DS32940
	DY(I+2) = DTEMP3	DS32950
	50 CONTINUE	DS32960
	RETURN	DS32970
	60 CONTINUE	DS32980
C		DS32990
C	CODE FOR EQUAL, POSITIVE, NONUNIT INCREMENTS.	DS33000
C		DS33010
	NS = N*INCX	DS33020
	DO 70 I=1,NS,INCX	DS33030
	DTEMP1 = DX(I)	DS33040
	DX(I) = DY(I)	DS33050
	DY(I) = DTEMP1	DS33060
	70 CONTINUE	DS33070
	RETURN	DS33080
	END	DS33090
	SUBROUTINE CSWAP(N,CX,INCX,CY,INCY)	CS33100
C		
C	INTERCHANGE COMPLEX CX AND COMPLEX CY	
C	FOR I = 0 TO N-1, INTERCHANGE CX(LX+I*INCX) AND CY(LY+I*INCY),	
C	WHERE LX = 1 IF INCX .GT. 0, ELSE LX = (-INCX)*N, AND LY IS	
C	DEFINED IN A SIMILAR WAY USING INCY.	
C		
	COMPLEX CX(1),CY(1),CTEMP	CS33140
C		CS33150
	IF(N .LE. 0)RETURN	CS33160
	IF(INCX.EQ.INCY.AND.INCX.GT.0) GO TO 20	CS33170
	KX = 1	CS33180
	KY = 1	CS33190
	IF(INCX.LT.0) KX = 1+(1-N)*INCX	CS33200
	IF(INCY.LT.0) KY = 1+(1-N)*INCY	CS33210
	DO 10 I = 1,N	CS33220
	CTEMP = CX(KX)	CS33230
	CX(KX) = CY(KY)	CS33240
	CY(KY) = CTEMP	CS33250
	KX = KX + INCX	CS33260
	KY = KY + INCY	CS33270
	10 CONTINUE	CS33280
	RETURN	CS33290
	20 CONTINUE	CS33300
	NS = N*INCX	CS33310
	DO 30 I=1,NS,INCX	CS33320
	CTEMP = CX(I)	CS33330
	CX(I) = CY(I)	CS33340
	CY(I) = CTEMP	CS33350
	30 CONTINUE	CS33360
	RETURN	CS33370
	END	CS33380


```

REAL FUNCTION SNRM2 ( N, SX, INCX)
INTEGER          NEXT
REAL            SX(1),  CUTLO, CUTHI, HITEST, SUM, XMAX, ZERO, ONE
DATA           ZERO, ONE /0.0E0, 1.0E0/

C
C   EUCLIDEAN NORM OF THE N-VECTOR STORED IN SX() WITH STORAGE
C   INCREMENT INCX .
C   IF N .LE. 0 RETURN WITH RESULT = 0.
C   IF N .GE. 1 THEN INCX MUST BE .GE. 1
C
C           C.L.LAWSON, 1978 JAN 08

C
C   FOUR PHASE METHOD      USING TWO BUILT-IN CONSTANTS THAT ARE
C   HOPEFULLY APPLICABLE TO ALL MACHINES.
C           CUTLO = MAXIMUM OF  SQRT(U/EPS)  OVER ALL KNOWN MACHINES.
C           CUTHI = MINIMUM OF  SQRT(V)      OVER ALL KNOWN MACHINES.
C   WHERE
C           EPS = SMALLEST NO. SUCH THAT EPS + 1. .GT. 1.
C           U   = SMALLEST POSITIVE NO.  (UNDERFLOW LIMIT)
C           V   = LARGEST NO.             (OVERFLOW LIMIT)
C
C   BRIEF OUTLINE OF ALGORITHM..
C
C   PHASE 1  SCANS ZERO COMPONENTS.
C   MOVE TO PHASE 2 WHEN A COMPONENT IS NONZERO AND .LE. CUTLO
C   MOVE TO PHASE 3 WHEN A COMPONENT IS .GT. CUTLO
C   MOVE TO PHASE 4 WHEN A COMPONENT IS .GE. CUTHI/M
C   WHERE M = N FOR X() REAL AND M = 2*N FOR COMPLEX.
C
C   VALUES FOR CUTLO AND CUTHI..
C   FROM THE ENVIRONMENTAL PARAMETERS LISTED IN THE IMSL CONVERTER
C   DOCUMENT THE LIMITING VALUES ARE AS FOLLOWS..
C   CUTLO, S.P.  U/EPS = 2**(-102) FOR HONEYWELL.  CLOSE SECONDS ARE
C               UNIVAC AND DEC AT 2**(-103)
C               THUS CUTLO = 2**(-51) = 4.44089E-16
C   CUTHI, S.P.  V = 2**127 FOR UNIVAC, HONEYWELL, AND DEC.
C               THUS CUTHI = 2**(63.5) = 1.30438E19
C   CUTLO, D.P.  U/EPS = 2**(-67) FOR HONEYWELL AND DEC.
C               THUS CUTLO = 2**(-33.5) = 8.23181D-11
C   CUTHI, D.P.  SAME AS S.P.  CUTHI = 1.30438D19
C   DATA CUTLO, CUTHI / 8.232D-11, 1.304D19 /
C   DATA CUTLO, CUTHI / 4.441E-16, 1.304E19 /
C   DATA CUTLO, CUTHI / 4.441E-16, 1.304E19 /

C
IF(N .GT. 0) GO TO 10
  SNRM2 = ZERO
  GO TO 300

C
10 ASSIGN 30 TO NEXT
  SUM = ZERO
  NN = N * INCX

C
C                               BEGIN MAIN LOOP

  I = 1
20  GO TO NEXT, (30, 50, 70, 110)
30  IF( ABS(SX(I)) .GT. CUTLO) GO TO 85
  ASSIGN 50 TO NEXT
  XMAX = ZERO

C
C                               PHASE 1.  SUM IS ZERO
C
50  IF( SX(I) .EQ. ZERO) GO TO 200
  DNRM2 = XMAX * DSQRT(SUM)
300 CONTINUE
  RETURN
  END

REAL FUNCTION SCNRM2( N, CX, INCX)
LOGICAL IMAG, SCALE
INTEGER          NEXT
REAL            CUTLO, CUTHI, HITEST, SUM, XMAX, ABSX, ZERO, ONE
COMPLEX        CX(1)
DATA           ZERO, ONE /0.0E0, 1.0E0/

```

```

C      UNITARY NORM OF THE COMPLEX N-VECTOR STORED IN CX() WITH STORAGE
C      INCREMENT INCX .
C      IF N .LE. 0 RETURN WITH RESULT = 0.
C      IF N .GE. 1 THEN INCX MUST BE .GE. 1
C
C      C.L.LAWSON , 1978 JAN 08
C
C      FOUR PHASE METHOD      USING TWO BUILT-IN CONSTANTS THAT ARE
C      HOPEFULLY APPLICABLE TO ALL MACHINES.
C      CUTLO = MAXIMUM OF  SQRT(U/EPS)  OVER ALL KNOWN MACHINES.
C      CUTHI = MINIMUM OF  SQRT(V)      OVER ALL KNOWN MACHINES.
C      WHERE
C      EPS = SMALLEST NO. SUCH THAT EPS + 1. .GT. 1.
C      U   = SMALLEST POSITIVE NO. (UNDERFLOW LIMIT)
C      V   = LARGEST NO. (OVERFLOW LIMIT)
C
C      BRIEF OUTLINE OF ALGORITHM..
C
C      PHASE 1  SCANS ZERO COMPONENTS.
C      MOVE TO PHASE 2 WHEN A COMPONENT IS NONZERO AND .LE. CUTLO
C      MOVE TO PHASE 3 WHEN A COMPONENT IS .GT. CUTLO
C      MOVE TO PHASE 4 WHEN A COMPONENT IS .GE. CUTHI/M
C      WHERE M = N FOR X() REAL AND M = 2*N FOR COMPLEX.
C
C      VALUES FOR CUTLO AND CUTHI..
C      FROM THE ENVIRONMENTAL PARAMETERS LISTED IN THE IMSL CONVERTER
C      DOCUMENT THE LIMITING VALUES ARE AS FOLLOWS..
C      CUTLO, S.P.  U/EPS = 2**(-102) FOR HONEYWELL.  CLOSE SECONDS ARE
C                  UNIVAC AND DEC AT 2**(-103)
C                  THUS CUTLO = 2**(-51) = 4.44089E-16
C      CUTHI, S.P.  V = 2**127 FOR UNIVAC, HONEYWELL, AND DEC.
C                  THUS CUTHI = 2**(63.5) = 1.30438E19
C      CUTLO, D.P.  U/EPS = 2**(-67) FOR HONEYWELL AND DEC.
C                  THUS CUTLO = 2**(-33.5) = 8.23181D-11
C      CUTHI, D.P.  SAME AS S.P.  CUTHI = 1.30438D19
C      DATA CUTLO, CUTHI / 8.232D-11, 1.304D19 /
C      DATA CUTLO, CUTHI / 4.441E-16, 1.304E19 /
C      DATA CUTLO, CUTHI / 4.441E-16, 1.304E19 /
C
C      IF(N .GT. 0) GO TO 10
C      SCNRM2 = ZERO
C      GO TO 300
C
C      10 ASSIGN 30 TO NEXT
C      SUM = ZERO
C      NN = N * INCX
C
C
C      BEGIN MAIN LOOP
C
C      DO 210 I=1,NN,INCX
C      ABSX = ABS(REAL(CX(I)))
C      IMAG = .FALSE.
C      GO TO NEXT,(30, 50, 70, 90, 110)
C      30 IF( ABSX .GT. CUTLO) GO TO 85
C      ASSIGN 50 TO NEXT
C      SCALE = .FALSE.
C
C
C      PHASE 1.  SUM IS ZERO
C
C      50 IF( ABSX .EQ. ZERO) GO TO 200
C      IF( ABSX .GT. CUTLO) GO TO 85
C
C
C      PREPARE FOR PHASE 2.
C
C      ASSIGN 70 TO NEXT
C      GO TO 105
C
C
C      PREPARE FOR PHASE 4.
C
C      100 ASSIGN 110 TO NEXT
C      SUM = (SUM / ABSX) / ABSX
C      105 SCALE = .TRUE.
C      XMAX = ABSX
C      GO TO 115
C
C
C      PHASE 2.  SUM IS SMALL.
C      SCALE TO AVOID DESTRUCTIVE UNDERFLOW.
C
C      70 IF( ABSX .GT. CUTLO ) GO TO 75

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C
C          COMMON CODE FOR PHASES 2 AND 4.
C          IN PHASE 4 SUM IS LARGE.  SCALE TO AVOID OVERFLOW.
C
110 IF( ABSX .LE. XMAX ) GO TO 115
    SUM = ONE + SUM * (XMAX / ABSX)**2
    XMAX = ABSX
    GO TO 200
C
115 SUM = SUM + (ABSX/XMAX)**2
    GO TO 200
C
C          PREPARE FOR PHASE 3.
C
75 SUM = (SUM * XMAX) * XMAX
C
85 ASSIGN 90 TO NEXT
    SCALE = .FALSE.
C
FOR REAL OR D.P. SET HITEST = CUTHI/N
FOR COMPLEX      SET HITEST = CUTHI/(2*N)
C
HITEST = CUTHI/FLOAT( N )
C
C          PHASE 3.  SUM IS MID-RANGE.  NO SCALING.
C
90 IF(ABSX .GE. HITEST) GO TO 100
    SUM = SUM + ABSX**2
200 CONTINUE
C          CONTROL SELECTION OF REAL AND IMAGINARY PARTS.
C
IF(IMAG) GO TO 210
    ABSX = ABS(AIMAG(CX(I)))
    IMAG = .TRUE.
    GO TO NEXT,( 50, 70, 90, 110 )
C
210 CONTINUE
    IMAG = .TRUE.
    GO TO NEXT,( 50, 70, 90, 110 )
C
210 CONTINUE
C
C          END OF MAIN LOOP.
C          COMPUTE SQUARE ROOT AND ADJUST FOR SCALING.
C
SCNRM2 = SQRT(SUM)
IF(SCALE) SCNRM2 = SCNRM2 * XMAX
300 CONTINUE
    RETURN
    END

REAL FUNCTION SASUM(N,SX,INCX)
C          SA35810
C          SA35820
C          RETURNS SUM OF MAGNITUDES OF SINGLE PRECISION SX.
C          SASUM = SUM FROM 0 TO N-1 OF ABS(SX(1+I*INCX))
C          SA35830
C          SA35840
REAL SX(1)
C          SA35850
SASUM = 0.0E0
C          SA35860
IF(N.LE.0)RETURN
C          SA35870
IF(INCX.EQ.1)GOTO 20
C          SA35880
C          SA35890
C          CODE FOR INCREMENTS NOT EQUAL TO 1.
C          SA35900
C          SA35910
NS = N*INCX
C          SA35920
DO 10 I=1,NS,INCX
C          SA35930
    SASUM = SASUM + ABS(SX(I))
C          SA35940
10 CONTINUE
C          SA35950
RETURN
C          SA35960
C          SA35970
C          CODE FOR INCREMENTS EQUAL TO 1.
C          SA35980
C          SA35990
C          SA36000
C          CLEAN-UP LOOP SO REMAINING VECTOR LENGTH IS A MULTIPLE OF 6.
C          SA36010
C          SA36020

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```

20 M = MOD(N,6)
   IF( M .EQ. 0 ) GO TO 40
   DO 30 I = 1,M
     SASUM = SASUM + ABS(SX(I))
30 CONTINUE
   IF( N .LT. 6 ) RETURN
40 MP1 = M + 1
   DO 50 I = MP1,N,6
     SASUM = SASUM + ABS(SX(I)) + ABS(SX(I + 1)) + ABS(SX(I + 2))
     $ + ABS(SX(I + 3)) + ABS(SX(I + 4)) + ABS(SX(I + 5))
50 CONTINUE
   RETURN
   END
SA36040
SA36050
SA36060
SA36070
SA36080
SA36090
SA36100
SA36110
SA36120
SA36130
SA36140
SA36150

DOUBLE PRECISION FUNCTION DASUM(N,DX,INCX)
C
C RETURNS SUM OF MAGNITUDES OF DOUBLE PRECISION DX.
C DASUM = SUM FROM 0 TO N-1 OF DABS(DX(1+I*INCX))
C
C DOUBLE PRECISION DX(1)
DASUM = 0.D0
IF(N.LE.0)RETURN
IF(INCX.EQ.1)GOTO 20
C
C CODE FOR INCREMENTS NOT EQUAL TO 1.
C
NS = N*INCX
DO 10 I=1,NS,INCX
  DASUM = DASUM + DABS(DX(I))
10 CONTINUE
RETURN
C
C CODE FOR INCREMENTS EQUAL TO 1.
C
C CLEAN-UP LOOP SO REMAINING VECTOR LENGTH IS A MULTIPLE OF 6.
C
20 M = MOD(N,6)
   IF( M .EQ. 0 ) GO TO 40
   DO 30 I = 1,M
     DASUM = DASUM + DABS(DX(I))
30 CONTINUE
   IF( N .LT. 6 ) RETURN
40 MP1 = M + 1
   DO 50 I = MP1,N,6
     DASUM = DASUM + DABS(DX(I)) + DABS(DX(I+1)) + DABS(DX(I+2))
     $ + DABS(DX(I+3)) + DABS(DX(I+4)) + DABS(DX(I+5))
50 CONTINUE
   RETURN
   END
DA36160
DA36170
DA36180
DA36190
DA36200
DA36210
DA36220
DA36230
DA36240
DA36250
DA36260
DA36270
DA36280
DA36290
DA36300
DA36310
DA36320
DA36330
DA36340
DA36350
DA36360
DA36370
DA36390
DA36400
DA36410
DA36420
DA36430
DA36440
DA36450
DA36460
DA36470
DA36480
DA36490
DA36500

FUNCTION SCASUM(N,CX,INCX)
C
C RETURNS SUMS OF MAGNITUDES OF REAL AND IMAGINARY PARTS OF
C COMPONENTS OF CX. NOTE THAT THIS IS NOT THE L1 NORM OF CX.
C CASUM = SUM FROM 0 TO N-1 OF ABS(REAL(CX(1+I*INCX))) +
C ABS(IMAG(CX(1+I*INCX)))
C
C COMPLEX CX(1)
C
SCASUM=0.
IF(N.LE.0) RETURN
NS = N*INCX
DO 10 I=1,NS,INCX
  SCASUM = SCASUM + ABS(REAL(CX(I))) + ABS(AIMAG(CX(I)))
10 CONTINUE
RETURN
END
SC36510
SC36560
SC36570
SC36580
SC36590
SC36600
SC36610
SC36620
SC36630
SC36640
SC36650

SUBROUTINE SSCAL(N,SA,SX,INCX)
C
C REPLACE SINGLE PRECISION SX BY SINGLE PRECISION SA*SX.
C FOR I = 0 TO N-1, REPLACE SX(1+I*INCX) WITH SA * SX(1+I*INCX)
SS36660
SS36670
SS36680

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C		SS36690
	REAL SA, SX(1)	SS36700
	IF(N.LE.0)RETURN	SS36710
	IF(INCX.EQ.1)GOTO 20	SS36720
C		SS36730
C	CODE FOR INCREMENTS NOT EQUAL TO 1.	SS36740
C		SS36750
	NS = N*INCX	SS36760
	DO 10 I = 1, NS, INCX	SS36770
	SX(I) = SA*SX(I)	SS36780
10	CONTINUE	SS36790
	RETURN	SS36800
C		SS36810
C	CODE FOR INCREMENTS EQUAL TO 1.	SS36820
C		SS36830
C		SS36840
C	CLEAN-UP LOOP SO REMAINING VECTOR LENGTH IS A MULTIPLE OF 5.	SS36850
C		SS36860
20	M = MOD(N,5)	
	IF(M .EQ. 0) GO TO 40	SS36880
	DO 30 I = 1, M	SS36890
	SX(I) = SA*SX(I)	SS36900
30	CONTINUE	SS36910
	IF(N .LT. 5) RETURN	SS36920
40	MP1 = M + 1	SS36930
	DO 50 I = MP1, N, 5	SS36940
	SX(I) = SA*SX(I)	SS36950
	SX(I + 1) = SA*SX(I + 1)	SS36960
	SX(I + 2) = SA*SX(I + 2)	SS36970
	SX(I + 3) = SA*SX(I + 3)	SS36980
	SX(I + 4) = SA*SX(I + 4)	SS36990
50	CONTINUE	SS37000
	RETURN	SS37010
	END	SS37020
	SUBROUTINE DSCAL(N,DA,DX,INCX)	DS37030
C		DS37040
C	REPLACE DOUBLE PRECISION DX BY DOUBLE PRECISION DA*DX.	DS37050
C	FOR I = 0 TO N-1, REPLACE DX(1+I*INCX) WITH DA * DX(1+I*INCX)	
C		DS37060
	DOUBLE PRECISION DA,DX(1)	DS37070
	IF(N.LE.0)RETURN	DS37080
	IF(INCX.EQ.1)GOTO 20	DS37090
C		DS37100
C	CODE FOR INCREMENTS NOT EQUAL TO 1.	DS37110
C		DS37120
	NS = N*INCX	DS37130
	DO 10 I = 1, NS, INCX	DS37140
	DX(I) = DA*DX(I)	DS37150
10	CONTINUE	DS37160
	RETURN	DS37170
C		DS37180
C	CODE FOR INCREMENTS EQUAL TO 1.	DS37190
C		DS37200
C		DS37210
C	CLEAN-UP LOOP SO REMAINING VECTOR LENGTH IS A MULTIPLE OF 5.	DS37220
C		DS37230
20	M = MOD(N,5)	
	IF(M .EQ. 0) GO TO 40	DS37250
	DO 30 I = 1, M	DS37260
	DX(I) = DA*DX(I)	DS37270
30	CONTINUE	DS37280
	IF(N .LT. 5) RETURN	DS37290
40	MP1 = M + 1	DS37300
	DO 50 I = MP1, N, 5	DS37310
	DX(I) = DA*DX(I)	DS37320
	DX(I + 1) = DA*DX(I + 1)	DS37330
	DX(I + 2) = DA*DX(I + 2)	DS37340
	DX(I + 3) = DA*DX(I + 3)	DS37350
	DX(I + 4) = DA*DX(I + 4)	DS37360
50	CONTINUE	DS37370
	RETURN	DS37380
	END	DS37390

```

SUBROUTINE CSCAL(N,CA,CX,INCX)
C
C REPLACE COMPLEX CX BY COMPLEX CA*CX.
C FOR I = 0 TO N-1, REPLACE CX(1+I*INCX) WITH CA * CX(1+I*INCX)
C
C COMPLEX CA,CX(1)
C
C IF(N .LE. 0) RETURN
NS = N*INCX
DO 10 I = 1,NS,INCX
CX(I) = CA*CX(I)
10 CONTINUE
RETURN
END
CS37400
CS37440
CS37450
CS37460
CS37470
CS37480
CS37490
CS37500
CS37510
CS37520

SUBROUTINE CSSCAL(N,SA,CX,INCX)
C
C REPLACE COMPLEX CX BY (SINGLE PRECISION SA) * (COMPLEX CX)
C FOR I = 0 TO N-1, REPLACE CX(1+I*INCX) WITH SA * CX(1+I*INCX)
C
C COMPLEX CX(1)
REAL SA
C
C IF(N .LE. 0) RETURN
NS = N*INCX
DO 10 I = 1,NS,INCX
CX(I) = SA*CX(I)
10 CONTINUE
RETURN
END
CS37530
CS37570
CS37580
CS37590
CS37600
CS37610
CS37620
CS37630
CS37640
CS37650
CS37660

INTEGER FUNCTION ISAMAX(N,SX,INCX)
C
C FIND SMALLEST INDEX OF MAXIMUM MAGNITUDE OF SINGLE PRECISION SX.
C ISAMAX = FIRST I, I = 1 TO N, TO MINIMIZE ABS(SX(1-INCX+I*INCX))
C
C REAL SX(1),SMAX,XMAG
ISAMAX = 0
IF(N.LE.0) RETURN
ISAMAX = 1
IF(N.LE.1) RETURN
IF(INCX.EQ.1) GOTO 20
C
C CODE FOR INCREMENTS NOT EQUAL TO 1.
C
C SMAX = ABS(SX(1))
NS = N*INCX
II = 1
DO 10 I=1,NS,INCX
XMAG = ABS(SX(I))
IF(XMAG.LE.SMAX) GO TO 5
ISAMAX = II
SMAX = XMAG
5 II = II + 1
10 CONTINUE
RETURN
C
C CODE FOR INCREMENTS EQUAL TO 1.
C
20 SMAX = ABS(SX(1))
DO 30 I = 2,N
XMAG = ABS(SX(I))
IF(XMAG.LE.SMAX) GO TO 30
ISAMAX = I
SMAX = XMAG
30 CONTINUE
RETURN
END
IS37670
IS37680
IS37690
IS37700
IS37710
IS37720
IS37730
IS37740
IS37750
IS37760
IS37770
IS37780
IS37790
IS37800
IS37810
IS37820
IS37830
IS37840
IS37850
IS37860
IS37870
IS37880
IS37890
IS37900
IS37910
IS37920
IS37930
IS37940
IS37950
IS37960
IS37970
IS37980
IS37990
IS38000
IS38010
IS38020

```

```

C          INTEGER FUNCTION IDAMAX(N,DX,INCX)                                ID38030
C                                                                                   ID38040
C          FIND SMALLEST INDEX OF MAXIMUM MAGNITUDE OF DOUBLE PRECISION DX.    ID38050
C          IDAMAX = FIRST I, I = 1 TO N, TO MINIMIZE ABS(DX(1-INCX+I*INCX))
C                                                                                   ID38060
C          DOUBLE PRECISION DX(1),DMAX,XMAG                                    ID38070
C          IDAMAX = 0                                                            ID38080
C          IF(N.LE.0) RETURN                                                    ID38090
C          IDAMAX = 1                                                            ID38100
C          IF(N.LE.1)RETURN                                                      ID38110
C          IF(INCX.EQ.1)GOTO 20                                                  ID38120
C                                                                                   ID38130
C          CODE FOR INCREMENTS NOT EQUAL TO 1.                                  ID38140
C                                                                                   ID38150
C          DMAX = DABS(DX(1))                                                    ID38160
C          NS = N*INCX                                                            ID38170
C          II = 1                                                                  ID38180
C              DO 10 I = 1,NS,INCX                                               ID38190
C                  XMAG = DABS(DX(I))                                           ID38200
C                  IF (XMAG.LE.DMAX) GO TO 5                                     ID38210
C                  IDAMAX = II                                                  ID38220
C                  DMAX = XMAG                                                  ID38230
C          5    II = II + 1                                                       ID38240
C          10   CONTINUE                                                         ID38250
C          RETURN                                                                  ID38260
C                                                                                   ID38270
C          CODE FOR INCREMENTS EQUAL TO 1.                                       ID38280
C                                                                                   ID38290
C          20   DMAX = DABS(DX(1))                                                ID38300
C              DO 30 I = 2,N                                                     ID38310
C                  XMAG = DABS(DX(I))                                           ID38320
C                  IF (XMAG.LE.DMAX) GO TO 30                                   ID38330
C                  IDAMAX = I                                                  ID38340
C                  DMAX = XMAG                                                  ID38350
C          30   CONTINUE                                                         ID38360
C          RETURN                                                                  ID38370
C          END                                                                    ID38380

C          INTEGER FUNCTION ICAMAX(N,CX,INCX)                                    IC38390
C                                                                                   IC38410
C          RETURNS THE INDEX OF THE COMPONENT OF CX HAVING THE                 IC38420
C          LARGEST SUM OF MAGNITUDES OF REAL AND IMAGINARY PARTS.
C          ICAMAX = FIRST I, I = 1 TO N, TO MINIMIZE
C          ABS(REAL(CX(1-INCX+I*INCX))) + ABS(IMAG(CX(1-INCX+I*INCX)))
C                                                                                   IC38440
C          COMPLEX CX(1)                                                         IC38450
C                                                                                   IC38460
C          ICAMAX = 0                                                             IC38470
C          IF(N.LE.0) RETURN                                                      IC38480
C          ICAMAX = 1                                                             IC38490
C          IF(N .LE. 1) RETURN                                                    IC38500
C          NS = N*INCX                                                            IC38510
C          II = 1                                                                  IC38520
C          SUMMAX = ABS(REAL(CX(1))) + ABS(AIMAG(CX(1)))                          IC38530
C              DO 20 I=1,NS,INCX                                                 IC38540
C                  SUMRI = ABS(REAL(CX(I))) + ABS(AIMAG(CX(I)))                IC38550
C                  IF(SUMMAX.GE.SUMRI) GO TO 10                                  IC38560
C                  SUMMAX = SUMRI                                               IC38570
C                  ICAMAX = II                                                  IC38580
C          10   II = II + 1                                                       IC38590
C          20   CONTINUE                                                         IC38600
C          RETURN                                                                  IC38610
C          END

C          BEGIN MAIN LOOP
C          I = 1
C          20   GO TO NEXT,(30, 50, 70, 110)
C          30   IF( DABS(DX(I)) .GT. CUTLO) GO TO 85
C          ASSIGN 50 TO NEXT
C          XMAX = ZERO

C          PHASE 1.  SUM IS ZERO
C          50   IF( DX(I) .EQ. ZERO) GO TO 2000
C              IF( DABS(DX(I)) .GT. CUTLO) GO TO 85

```

```

C
C
C          PREPARE FOR PHASE 2.
C          ASSIGN 70 TO NEXT
C          GO TO 105
C
C
C          PREPARE FOR PHASE 4.
C
C          100 I = J
C          ASSIGN 110 TO NEXT
C          SUM = (SUM / DX(I)) / DX(I)
C          105 XMAX = DABS(DX(I))
C          GO TO 115
C
C          PHASE 2.  SUM IS SMALL.
C          SCALE TO AVOID DESTRUCTIVE UNDERFLOW.
C
C          70 IF( DABS(DX(I)) .GT. CUTLO ) GO TO 75
C
C          COMMON CODE FOR PHASES 2 AND 4.
C          IN PHASE 4 SUM IS LARGE.  SCALE TO AVOID OVERFLOW.
C
C          110 IF( DABS(DX(I)) .LE. XMAX ) GO TO 115
C          SUM = ONE + SUM * (XMAX / DX(I))**2
C          XMAX = DABS(DX(I))
C          GO TO 200
C
C          115 SUM = SUM + (DX(I)/XMAX)**2
C
C          GO TO 200
C
C
C          PREPARE FOR PHASE 3.
C
C          75 SUM = (SUM * XMAX) * XMAX
C
C          FOR REAL OR D.P. SET HITEST = CUTHI/N
C          FOR COMPLEX      SET HITEST = CUTHI/(2*N)
C
C          85 HITEST = CUTHI/FLOAT( N )
C
C          PHASE 3.  SUM IS MID-RANGE.  NO SCALING.
C
C          DO 95 J =I,NN,INCX
C          IF(DABS(DX(J)) .GE. HITEST) GO TO 100
C          95  SUM = SUM + DX(J)**2
C          DNRM2 = DSQRT( SUM )
C          GO TO 300
C
C          200 CONTINUE
C          I = I + INCX
C          IF ( I .LE. NN ) GO TO 20
C
C          END OF MAIN LOOP.
C
C          COMPUTE SQUARE ROOT AND ADJUST FOR SCALING.
C
C          I = I + INCX
C          IF ( I .LE. NN ) GO TO 20
C
C          END OF MAIN LOOP.
C
C          COMPUTE SQUARE ROOT AND ADJUST FOR SCALING.
C
C          SNRM2 = XMAX * SQRT(SUM)
C          300 CONTINUE
C          RETURN
C          END

```

DOUBLE PRECISION FUNCTION DNRM2 (N, DX, INCX)
 INTEGER NEXT
 DOUBLE PRECISION DX(1), CUTLO, CUTHI, HITEST, SUM, XMAX,ZERO,ONE
 DATA ZERO, ONE /0.0D0, 1.0D0/

```

C
C          EUCLIDEAN NORM OF THE N-VECTOR STORED IN DX() WITH STORAGE

```



```

C      INCREMENT INCX .
C      IF N .LE. 0 RETURN WITH RESULT = 0.
C      IF N .GE. 1 THEN INCX MUST BE .GE. 1
C
C      C.L.LAWSON, 1978 JAN 08
C
C      FOUR PHASE METHOD      USING TWO BUILT-IN CONSTANTS THAT ARE
C      HOPEFULLY APPLICABLE TO ALL MACHINES.
C      CUTLO = MAXIMUM OF DSQRT(U/EPS) OVER ALL KNOWN MACHINES.
C      CUTHI = MINIMUM OF DSQRT(V) OVER ALL KNOWN MACHINES.
C      WHERE
C      EPS = SMALLEST NO. SUCH THAT EPS + 1. .GT. 1.
C      U = SMALLEST POSITIVE NO. (UNDERFLOW LIMIT)
C      V = LARGEST NO. (OVERFLOW LIMIT)
C
C      BRIEF OUTLINE OF ALGORITHM..
C
C      PHASE 1 SCANS ZERO COMPONENTS.
C      MOVE TO PHASE 2 WHEN A COMPONENT IS NONZERO AND .LE. CUTLO
C      MOVE TO PHASE 3 WHEN A COMPONENT IS .GT. CUTLO
C      MOVE TO PHASE 4 WHEN A COMPONENT IS .GE. CUTHI/M
C      WHERE M = N FOR X() REAL AND M = 2*N FOR COMPLEX.
C
C      VALUES FOR CUTLO AND CUTHI..
C      FROM THE ENVIRONMENTAL PARAMETERS LISTED IN THE IMSL CONVERTER
C      DOCUMENT THE LIMITING VALUES ARE AS FOLLOWS..
C      CUTLO, S.P. U/EPS = 2**(-102) FOR HONEYWELL. CLOSE SECONDS ARE
C      UNIVAC AND DEC AT 2**(-103)
C      THUS CUTLO = 2**(-51) = 4.44089E-16
C      CUTHI, S.P. V = 2**127 FOR UNIVAC, HONEYWELL, AND DEC.
C      THUS CUTHI = 2**(63.5) = 1.30438E19
C      CUTLO, D.P. U/EPS = 2**(-67) FOR HONEYWELL AND DEC.
C      THUS CUTLO = 2**(-33.5) = 8.23181D-11
C
C      CUTHI, D.P. SAME AS S.P. CUTHI = 1.30438D19
C      DATA CUTLO, CUTHI / 8.232D-11, 1.304D19 /
C      DATA CUTLO, CUTHI / 4.441E-16, 1.304E19 /
C      DATA CUTLO, CUTHI / 8.232D-11, 1.304D19 /
C
C      IF(N .GT. 0) GO TO 10
C      DNRM2 = ZERO
C      GO TO 300
C
C      10 ASSIGN 30 TO NEXT
C      SUM = ZERO
C      NN = N * INCX
C
C      DOUBLE PRECISION FUNCTION DQDOTI(N,DB,QC,DX,INCX,DY,INCY)
C      D.P. DOT PRODUCT WITH EXTENDED PRECISION ACCUMULATION (AND RESULT)
C      QC AND DQDOTI ARE SET = DB + SUM FOR I = 0 TO N-1 OF
C      DX(LX+I*INCX) * DY(LY+I*INCY), WHERE QC IS AN EXTENDED
C      PRECISION RESULT WHICH CAN BE USED AS INPUT TO DQDOTA,
C      AND LX = 1 IF INCX .GE. 0, ELSE LX = (-INCX)*N, AND LY IS
C      DEFINED IN A SIMILAR WAY USING INCY. THE MP PACKAGE BY
C      RICHARD P. BRENT IS USED FOR THE EXTENDED PRECISION ARITHMETIC.
C
C      FRED T. KROGH, JPL, 1977, JUNE 1
C2
C      DOUBLE PRECISION DX(1), DY(1), DB
C      INTEGER QC(10), QX(10), QY(10)
C      THE COMMON BLOCK FOR THE MP PACKAGE (MODIFIED TO GIVE IT A NAME)
C      COMMON /MPCOM/ MPB, MPT, MPM, MPLUN, MPMXR, MPR(12)
C      DATA I1 / 0 /
C      IF I1 IS 0 THE MP PACKAGE MUST BE INITIALIZED (MPBLAS SETS I1 = 1)
C      IF (I1 .EQ. 0) CALL MPBLAS(I1)
C      QC(1) = 0
C      IF (DB .EQ. 0.D0) GO TO 60
C      CALL MPCDM(DB, QX)
C      CALL MPADD(QC, QX, QC)
C      60 IF (N .EQ. 0) GO TO 80
C      IX = 1
C      IY = 1
C      IF (INCX .LT. 0) IX = (-N + 1) * INCX + 1
C      IF (INCY .LT. 0) IY = (-N + 1) * INCY + 1
C      DO 70 I = 1, N

```

```

      CALL MPCDM(DX(IX), QX)
      CALL MPCDM(DY(IY), QY)
      CALL MPMUL(QX, QY, QX)
      CALL MPADD(QC, QX, QC)
      IX = IX + INCX
      IY = IY + INCY
70 CONTINUE
80 CALL MPCMD(QC, DQDOTI)
   RETURN
   END

      SUBROUTINE DROTM (N,DX,INCX,DY,INCY,DPARAM)
C
C   APPLY THE MODIFIED GIVENS TRANSFORMATION, H, TO THE 2 BY N MATRIX
C   (DX**T) , WHERE **T INDICATES TRANSPOSE. THE ELEMENTS OF DX ARE IN
C   (DY**T)
C
C   DX(LX+I*INCX), I = 0 TO N-1, WHERE LX = 1 IF INCX .GE. 0, ELSE
C   LX = (-INCX)*N, AND SIMILARLY FOR SY USING LY AND INCY.
C   WITH DPARAM(1)=DFLAG, H HAS ONE OF THE FOLLOWING FORMS..
C
C   DFLAG=-1.D0      DFLAG=0.D0      DFLAG=1.D0      DFLAG=-2.D0
C
C   (DH11 DH12)      (1.D0 DH12)      (DH11 1.D0)      (1.D0 0.D0)
C   H=( )            ( )              ( )              ( )
C   (DH21 DH22),      (DH21 1.D0),      (-1.D0 DH22),      (0.D0 1.D0).
C   SEE DROTMG FOR A DESCRIPTION OF DATA STORAGE IN DPARAM.
C
C   DOUBLE PRECISION DFLAG,DH12,DH22,DX,TWO,Z,DH11,DH21,
1 DPARAM,DY,W,ZERO
   DIMENSION DX(1),DY(1),DPARAM(5)
   DATA ZERO,TWO/0.D0,2.D0/
C
   DFLAG=DPARAM(1)
   IF(N .LE. 0 .OR. (DFLAG+TWO.EQ.ZERO)) GO TO 140
   IF(.NOT.(INCX.EQ.INCY.AND. INCX .GT.0)) GO TO 70
C
      NSTEPS=N*INCX
      IF(DFLAG) 50,10,30
10 CONTINUE
      DH12=DPARAM(4)
      DH21=DPARAM(3)
         DO 20 I=1,NSTEPS,INCX
            W=DX(I)
            Z=DY(I)
            DX(I)=W+Z*DH12
            DY(I)=W*DH21+Z
20 CONTINUE
      GO TO 140
30 CONTINUE
      DH11=DPARAM(2)
      DH22=DPARAM(5)
         DO 40 I=1,NSTEPS,INCX
            W=DX(I)
            Z=DY(I)
            DX(I)=W*DH11+Z
            DY(I)=-W+DH22*Z
40 CONTINUE
      GO TO 140
50 CONTINUE
      DH11=DPARAM(2)
      DH12=DPARAM(4)
      DH21=DPARAM(3)
      DH22=DPARAM(5)
         DO 60 I=1,NSTEPS,INCX
            W=DX(I)
            Z=DY(I)
            DX(I)=W*DH11+Z*DH12
            DY(I)=W*DH21+Z*DH22
60 CONTINUE
      GO TO 140
70 CONTINUE
      KX=1

```

	KY=1	DR30080
	IF (INCX .LT. 0) KX=1+(1-N)*INCX	DR30090
	IF (INCY .LT. 0) KY=1+(1-N)*INCY	DR30100
C		DR30110
	IF (DFLAG) 120, 80, 100	DR30120
80	CONTINUE	DR30130
	DH12=DPARAM(4)	DR30140
	DH21=DPARAM(3)	DR30150
	DO 90 I=1, N	DR30160
	W=DX (KX)	DR30170
	Z=DY (KY)	DR30180
	DX (KX)=W+Z*DH12	DR30190
	DY (KY)=W*DH21+Z	DR30200
	KX=KX+INCX	DR30210
	KY=KY+INCY	DR30220
90	CONTINUE	DR30230
	GO TO 140	DR30240
100	CONTINUE	DR30250
	DH11=DPARAM(2)	DR30260
	DH22=DPARAM(5)	DR30270
	DO 110 I=1, N	DR30280
	W=DX (KX)	DR30290
	Z=DY (KY)	DR30300
	DX (KX)=W*DH11+Z	DR30310
	DY (KY)=-W+DH22*Z	DR30320
	KX=KX+INCX	DR30330
	KY=KY+INCY	DR30340
110	CONTINUE	DR30350
	GO TO 140	DR30360
120	CONTINUE	DR30370
	DH11=DPARAM(2)	DR30380
	DH12=DPARAM(4)	DR30390
	DH21=DPARAM(3)	DR30400
	DH22=DPARAM(5)	DR30410
	DO 130 I=1, N	DR30420
	W=DX (KX)	DR30430
	Z=DY (KY)	DR30440
	DX (KX)=W*DH11+Z*DH12	DR30450
	DY (KY)=W*DH21+Z*DH22	DR30460
	KX=KX+INCX	DR30470
	KY=KY+INCY	DR30480
130	CONTINUE	DR30490
140	CONTINUE	DR30500
	RETURN	DR30510
	END	DR30520

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REMARK ON ALGORITHM 539

Basic Linear Algebra Subprograms for Fortran Usage [C.L. Lawson, R.J. Hanson, D.R. Kincaid, and F.T. Krogh, *ACM Trans. Math. Softw.* 5, 3 (Sept. 1979), 324-325]

David S. Dodson and Roger G. Grimes [Received 11 May 1982; revised 25 August 1982; accepted 17 September 1982]

Boeing Computer Services Company, Mail Stop 9C-01, 565 Andover Park West, Tukwila, WA 98188.

The companion [5] to Algorithm 539 (the BLAS) contains two errors which we discovered in the preparation of [1].

The more serious error occurs in the mathematical specifications for subroutines **SROTG** and **DROTG**, which construct Givens plane rotations. For convenience of reference, we include the relevant specifications from [5].

Given a and b , each of these subroutines computes

$$\sigma = \begin{cases} \text{sgn}(a) & \text{if } |a| > |b|, \\ \text{sgn}(b) & \text{if } |b| \geq |a|, \end{cases} \quad r = \sigma(a^2 + b^2)^{1/2}, \quad (1a,b)$$

Table I. Status of BLAS Implementations

Source	SROTG and DROTG		
	Reference	Documentation	Code
Boeing Computer Services	[1]	Correct	Correct
Dongarra et al.	[2]	Incorrect	Incorrect
Floating Point Systems	[3]	Incorrect	Incorrect
IMSL	[4]	Incorrect	Correct
Lawson et al.	[5]	Incorrect	Correct
Petersen	[6]	Correct	Incorrect

$$= \begin{cases} a/r & \text{if } r \neq 0, \\ 1 & \text{if } r = 0, \end{cases} \quad s = \begin{cases} b/r & \text{if } r \neq 0, \\ 1 & \text{if } r = 0, \end{cases} \quad (1c,d)$$

$$z = \begin{cases} s & \text{if } |a| > |b|, \\ 1/c & \text{if } |b| \geq |a| \text{ and } c \neq 0, \\ 1 & \text{if } c = 0. \end{cases} \quad (1e)$$

If the user later wishes to reconstruct c and s from z , it can be done as follows:

If $z = 1$ set $c = 0$ and $s = 1$.

If $|z| < 1$ set $c = (1 - z^2)^{1/2}$ and $s = z$. (2)

If $|z| > 1$ set $c = 1/z$ and $s = (1 - c^2)^{1/2}$.

The problem occurs in the computation of z . In particular, when $a = b = 0$, (1a-e) yield the results $r = 0$, $c = 1$, $s = 0$, and $z = 1$. However, when c and s are reconstructed from z using (2), the incorrect values $c = 0$ and $s = 1$ result. This discrepancy can be resolved only by changing the computation of z , as $z = 1$ can result both when $c = 0$ (when $a = 0$ and $b \neq 0$) and when $c = 1$ (when $a = b = 0$).

Stewart [7] intended that z represent the smaller in magnitude of c and s and that it indicate which one of c and s is the smaller, since the magnitude of the larger can be reconstructed stably from the smaller. He expressly omitted consideration of the point $a = b = 0$, which is the only point at which (1e) fails. We propose that z be computed directly from c and s . Two observations simplify matters. First, the larger in magnitude of c and s is positive: $|s| < |c|$ implies $c > 0$ and $0 < |c| \leq |s|$ implies $s > 0$. Second, if $c = 0$, then $s = 1$. Thus the value z can be redefined as follows:

$$z = \begin{cases} s & \text{if } |s| < c \text{ or } c = 0, \\ 1/c & \text{if } 0 < |c| \leq s. \end{cases} \quad (3)$$

This definition of z in (3) differs from (1a-e) only at the single point $a = b = 0$, where $z = 0$ results. For $z = 0$, the reconstruction formulas (2) give the correct values, $c = 1$ and $s = 0$.

Even though the original mathematical specifications for subroutines **SROTG** and **DROTG** are incorrect, the FORTRAN and assembly language versions of **SROTG** and **DROTG**, available from the ACM Algorithm Distribution Service as part of Algorithm 539, do conform to the correct mathematical specification, as given in this Remark.

BLAS documentation has appeared in several publications, and BLAS code is available for use or distribution from several sources. Table I lists the status of some of these BLAS implementations.

A second error, more innocuous than the above, involves the dimensions of the arrays passed to the BLAS. In several places, the length is incorrectly specified as $\max(1, N * |\text{INCX}|)$. The correct length is $\max(1, 1 + (N - 1) * |\text{INCX}|)$. This error occurs near the beginnings of Sections 5 and 7, and would be unobtrusive if Section 7 did not contain the word "precisely."

ACKNOWLEDGMENT

We thank Dr. C. L. Lawson for his helpful comments, which have led to a redefinition of z that is superior to what we had originally suggested.

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ALGORITHM 540

PDECOL, General Collocation Software for Partial Differential Equations [D3]

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Key Words and Phrases: collocation methods, partial differential equations, numerical software, method of lines

CR Categories: 3.20, 3.22, 4.0, 5.17

Language: Fortran

DESCRIPTION

1. Introduction

The basic purpose of this paper is to describe and discuss PDECOL, which is a new computer software package for numerically solving coupled systems of nonlinear partial differential equations (PDE's) in one space and one time dimension. The package implements finite element collocation methods based on piecewise polynomials for the spatial discretization techniques. The time integration process is then accomplished by widely acceptable procedures [7] which are generalizations of the usual methods for treating time dependent partial differential equations.

PDECOL is unique because of its flexibility both in the class of problems it addresses and in the variety of methods it provides for use in the solution process. High order methods (as well as low order ones) are readily available for use in both the spatial and time discretization procedures. The time integration methods used feature automatic time step size and integration formula order selection so as to efficiently solve the problem at hand and yet achieve a user specified time integration error level.

PDECOL consists of a collection of 19 subroutines written in reasonably standard Fortran, and therefore is quite portable and can be readily used on almost any capable computer with a Fortran compiler. No special hardware features are required.

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PDECOL is designed to solve broad classes of difficult systems of partial differential equations that describe physical processes. This package should be of interest to almost anyone involved in scientific and engineering simulations, calculations, and model development. It should also be of interest to those involved in mathematical software development and, in particular, to those concerned with software development for partial differential equations.

One of our main objectives in producing PDECOL was to make available a package which will allow users to try out a variety of collocation techniques on a variety of problems with little program development required. As is mentioned in the text, there are over 3000 different potentially useful PDE methods in PDECOL. Basically, we are relying on others to produce the final judgments as to the overall effectiveness of collocation for their problems. We feel from our experience that collocation looks attractive enough to warrant serious consideration by others—hence, the package PDECOL.

2. Class of Problems

PDECOL is designed to solve the general system of *NPDE* nonlinear partial differential equations of at most second order on the interval $[x_L, x_R]$ for $t \geq t_0$ which is of the form

$$\frac{\partial \mathbf{u}}{\partial t} = \mathbf{f}(t, x, \mathbf{u}, \mathbf{u}_x, \mathbf{u}_{xx}), \quad (2.1)$$

where

$$\begin{aligned} \mathbf{u} &= (u_1, u_2, \dots, u_{NPDE}), \\ \mathbf{u}_x &= \left(\frac{\partial u_1}{\partial x}, \frac{\partial u_2}{\partial x}, \dots, \frac{\partial u_{NPDE}}{\partial x} \right), \\ \mathbf{u}_{xx} &= \left(\frac{\partial^2 u_1}{\partial x^2}, \frac{\partial^2 u_2}{\partial x^2}, \dots, \frac{\partial^2 u_{NPDE}}{\partial x^2} \right). \end{aligned} \quad (2.2)$$

Each u_k is a function of the scalar quantities t and x , $k = 1, 2, \dots, NPDE$. In (2.1) \mathbf{f} represents an arbitrary vector valued function whose *NPDE* components define the respective partial differential equations of the PDE system. As there is no explicit requirement that \mathbf{f} actually depends on \mathbf{u}_x and/or \mathbf{u}_{xx} , any particular equation in (2.1) may actually be an ordinary differential equation (ODE), a first-order PDE, or a second-order PDE.

Depending on the particular type of equation, 0, 1, or 2 boundary conditions may be required for each equation in the system (2.1). These are imposed at x_L and/or x_R (or not imposed at all in the case of no condition) and must be of the form

$$\mathbf{b}(\mathbf{u}, \mathbf{u}_x) = \mathbf{z}(t), \quad (2.3)$$

where \mathbf{b} and \mathbf{z} are arbitrary vector valued functions with *NPDE* components and \mathbf{u} , \mathbf{u}_x , and t are as above. We make the basic assumption that these boundary conditions (2.3) must be consistent with the initial conditions (2.4) which are described next.

Each solution component u_k is assumed to be a known function of x at the initial time $t = t_0$. That is,

$$u_k(t_0, x) = \phi_k(x), \quad k = 1, 2, \dots, NPDE, \quad (2.4)$$

where each $\phi_k(x)$ is a known function of x . The initial condition functions must be consistent with the boundary conditions (2.3), (i.e. the initial condition functions must satisfy the boundary conditions for $t = t_0$).

We assume that all functions are continuous in t and at least piecewise continuous in x . With some understanding of the internal workings of the package, it is possible to solve problems with inconsistent boundary and initial conditions

or problems with jump discontinuities in the boundary conditions with respect to t .

It is obvious that the possible PDE problems described above are sufficiently general to include many systems for which solutions may not exist or may not be unique. *It must be, of course, the user's responsibility to define a mathematically meaningful PDE problem.*

3. Defining the PDE Problem

The PDE system described above is completely specified if one defines: *NPDE*; the interval $[x_L, x_R]$; the initial time t_0 ; the vector functions \mathbf{f} , \mathbf{b} , and \mathbf{z} ; and the initial condition functions $\phi_k(x)$, $k = 1, 2, \dots, NPDE$. The PDE problem can be properly defined for solution by PDECOL by constructing a main program and three subprograms, F, BNDRY, and UINIT, which specify the above-mentioned quantities.

Main Program. The construction of the main program that a user is required to supply is usually straightforward. The main program serves four basic purposes. First, initialization of the calling arguments for PDECOL must be performed. Second, the main program must call the package. Third, it should provide logic to detect possible error returns from PDECOL. Fourth, the main program should perform the desired output of the computed results. A typical main program is shown in a later section for one numerical example.

Required User Supplied Subroutines. The user is required to construct three subroutines which define the form of the PDE problem. Stated briefly, the purposes of these routines, F, BNDRY, and UINIT, are as follows. For given input values of t and x and corresponding input values of \mathbf{u} , \mathbf{u}_x , and \mathbf{u}_{xx} which are associated with this time and spatial position: subroutine F is to compute appropriate values for the functions f_k in (2.1); subroutine BNDRY is to compute appropriate values for the derivatives of the boundary condition functions b_k and z_k in (2.3) at the left or right boundary as determined by the value of x (see (4.3) and (4.4)); and subroutine UINIT is to compute initial condition function values (2.4).

The subroutines F, BNDRY, and UINIT are usually easily constructed and an example showing these routines is contained in a later section. More specific comments are given in the Algorithm at the end of this paper.

4. Methods Used

The software package PDECOL is based on the method of lines [7, 8, 11] and uses a finite element collocation procedure (with piecewise polynomials as the trial space) for the discretization of the spatial variable x . The collocation procedure reduces the PDE system to a semidiscrete system (actually an initial-value ODE system), which then depends only on the time variable t . The time integration is then accomplished by use of slightly modified standard techniques [5, 6, 7], which will be discussed briefly in a later section.

Piecewise Polynomials. The user has the opportunity to specify the piecewise polynomial space which is to be used to compute his approximate solution. In selecting this space the order, *KORD*, of the polynomials to be used must first be specified (*KORD* = polynomial degree + 1). Next, the number of pieces (intervals), *NINT*, into which the spatial domain $[x_L, x_R]$ is to be divided is chosen. The *NINT* + 1 distinct breakpoints of the domain must be defined and set into the array *XBKPT* in strictly increasing order, i.e.

$$x_L = XBKPT(1) < XBKPT(2) < \dots < XBKPT(NINT + 1) = x_R.$$

The approximate solution at any time, t , will be a polynomial of order *KORD* in each subinterval $[XBKPT(i), XBKPT(i + 1)]$, $i = 1, 2, \dots, NINT$. The number of continuity conditions, *NCC*, to be imposed on the polynomial pieces across all of the interior breakpoints is the last piece of user supplied data which is required

to uniquely determine the desired piecewise polynomial space. For example, $NCC = 2$ would require that the approximate solution (made up of the separate polynomial pieces) and its first spatial derivative be continuous at the breakpoints and hence on the entire domain $[x_L, x_R]$. $NCC = 3$ would require in addition that the second spatial derivative be continuous. The dimension of this linear space is known and finite and is $NCPTS = KORD * NINT - NCC * (NINT - 1)$. The well-known B-spline basis [2] for this space is used by PDECOL and it consists of $NCPTS$ known piecewise polynomial functions $\Phi_i(x)$, $i = 1, 2, \dots, NCPTS$, which do not depend on the time variable t . Use of the B-spline basis results in banded (in contrast to full) matrix problems to be solved in the package. The computer program requires $3 \leq KORD \leq 20$, $1 < NCC < KORD$, and $NINT \geq 1$.

Collocation over Piecewise Polynomials. The basic assumption made is that at any given time t , each approximate solution component, u_k , is a piecewise polynomial in the user specified space and hence can be written in terms of the B-spline basis functions as

$$u_k(t, x) = \sum_{i=1}^{NCPTS} c_{i,k}(t) \Phi_i(x), \quad k = 1, 2, \dots, NPDE. \quad (4.1)$$

The unknown coefficients $c_{i,k}$ depend only on the time t , and the *known* basis functions Φ_i depend on x . The semidiscrete equations (actually ordinary differential equations) which determine these coefficients, $c_{i,k}$, for $i = 1, 2, \dots, NCPTS$ and $k = 1, 2, \dots, NPDE$, are obtained by collocating, i.e. by requiring the approximate $u_k(t, x)$ in (4.1) to satisfy the PDE's (2.1) and the boundary conditions (2.3) exactly, at a set of $NCPTS$ collocation points.

To be more specific, we choose $NCPTS$ collocation points such that

$$x_L = \xi_1 < \xi_2 < \dots < \xi_{NCPTS} = x_R$$

and $\Phi_i(\xi_i) \neq 0$ for $i = 1, 2, \dots, NCPTS$. Then, substituting (4.1) into (2.1) and requiring (2.1) to be valid at the interior collocation points gives

$$\sum_{i=1}^{NCPTS} \Phi_i(\xi_j) \frac{dc_{i,k}}{dt} = f_k(t, \xi_j, \mathbf{u}(t, \xi_j), \mathbf{u}_x(t, \xi_j), \mathbf{u}_{xx}(t, \xi_j)),$$

$$j = 2, 3, \dots, NCPTS - 1, \quad k = 1, 2, \dots, NPDE. \quad (4.2)$$

To determine the equations corresponding to $j = 1$ and $j = NCPTS$, we form equations which depend on the type of boundary condition. It will suffice to show the technique we use for the left boundary, $x = x_L$ where $j = 1$, since those for the right boundary, $x = x_R$ where $j = NCPTS$, are completely analogous.

Normally, we form an ODE corresponding to the point $x = x_L$ by differentiating the boundary conditions (2.3) with respect to t , which gives

$$\sum_{j=1}^{NPDE} \left\{ \frac{\partial b_k}{\partial u_j} \frac{\partial u_j}{\partial t} + \frac{\partial b_k}{\partial u_{x_j}} \frac{\partial u_{x_j}}{\partial t} \right\} = \frac{dz_k}{dt}. \quad (4.3)$$

Substituting (4.1) into (4.3) and using the facts [1, 3] that $\Phi_1(x_L) \neq 0$, $\Phi_i(x_L) = 0$, $i = 2, 3, \dots, NCPTS$, and that $\Phi'_1(x_L) \neq 0$, $\Phi'_2(x_L) \neq 0$, $\Phi'_i(x_L) = 0$, $i = 3, 4, \dots, NCPTS$ gives the appropriate ODE:

$$\sum_{j=1}^{NPDE} \left\{ \frac{\partial b_k}{\partial u_j} \Phi_1(x_L) + \frac{\partial b_k}{\partial u_{x_j}} \Phi'_1(x_L) \right\} \frac{dc_{1,j}}{dt} + \sum_{j=1}^{NPDE} \left\{ \frac{\partial b_k}{\partial u_{x_j}} \Phi'_2(x_L) \right\} \frac{dc_{2,j}}{dt} = \frac{dz_k}{dt}. \quad (4.4)$$

In the special case when no boundary condition is desired for the $k = k_0$ equation, we simply collocate in the usual manner at the boundary point and obtain (4.2) with $j = 1$ and $k = k_0$ for our equation corresponding to the point $x = x_L$.

Combining the boundary condition equations with (4.2) yields a semidiscrete system of $N = NPDE * NCPTS$ time dependent ordinary differential equations which have the form

$$A \frac{d\mathbf{c}}{dt} = \mathbf{g}(t, \mathbf{c}). \quad (4.5)$$

The matrix A in (4.5) and the Jacobian matrix of \mathbf{g} , $\partial\mathbf{g}/\partial\mathbf{c}$, are matrices with a maximum bandwidth of $2 * (KORD - 1) * NPDE - 1$ except when $KORD = 3$ and an equation with no boundary condition at one endpoint exists. Except for the first and last block rows, all of the entries of A are simply basis function values at the collocation points. The first and last block rows of A consist of the appropriate boundary condition equation coefficients from (4.4) or (4.2). Since at any point at most $KORD$ basis functions have nonzero values, each block row of A consists of $KORD$ $NPDE$ by $NPDE$ matrices.

We remind the reader that the unknowns in (4.5) which are actually computed by PDECOL are the *basis function coefficients* in (4.1) and *not* the actual approximate solution values. However, with a knowledge of these coefficients, we can then easily evaluate the approximate solution values by using (4.1).

Accuracy Considerations. There are two sources of error in the approximate solutions generated by PDECOL. The first is due to the time discretization methods used and the second is due to the collocation spatial discretization technique. PDECOL attempts to control the time discretization error and to maintain it below a user specified level by dynamically selecting appropriate time step sizes and time integration formulas. We refer the reader to the package documentation for the details on how this is accomplished. Control of the error introduced by the spatial discretization is a more difficult problem, and we make the following observations.

Piecewise polynomials of order $KORD$ have the approximation property [13] that sufficiently smooth functions can be approximated by these polynomials such that the spatial errors in the approximation are proportional to h^{KORD} , where $h = \max_i [XBKPT(i + 1) - XBKPT(i)]$. Derivatives of order j of these smooth functions are also approximated with spatial errors which are proportional to h^{KORD-j} . We can use these approximation properties to estimate the expected orders of accuracy (for the *spatial* discretization) of the collocation techniques implemented in PDECOL. In particular, when using a piecewise polynomial space of order $KORD$, since eq. (2.1) involves second-order spatial derivatives, we expect PDECOL to generate approximate solutions with *spatial* errors which are proportional to h^{KORD-2} . However, for the special spaces with $KORD > 3$ and $NCC = 2$, a special choice of collocation points (Gauss-Legendre quadrature points in each subinterval) generates approximate solutions with *spatial* errors which are proportional to h^{KORD} for certain classes of PDE problems [4].

We emphasize that the piecewise polynomial space used in PDECOL (*which is selected by the user*) will determine the magnitude of the spatial discretization errors in the computed approximate solution. *The package has no control over errors introduced by the user's choice of the piecewise polynomial space.*

5. Limitations and Use of PDECOL

Of course, it goes without saying that no one program such as PDECOL will solve all PDE problems. There are several reasons for this, and we will enumerate them.

First, in order to be "solvable" a PDE problem should be properly posed in a mathematical sense. There is nothing that PDECOL can do to insure or detect that a user's problem is properly posed and so we reiterate that this responsibility must fall back to the user. Discussion of what constitutes a properly posed problem is beyond the scope and intent of this paper and we refer the interested reader to any good PDE textbook.

Second, PDECOL (and almost any general purpose software package) imposes certain restrictions on the classes of problems it allows. These restrictions are enumerated in Section 2. Perhaps the most unpleasant of these restrictions is the form of the boundary conditions (2.3) and the fact that the boundary conditions must be consistent with the initial conditions. Maintaining high order approxi-

mations at the boundary is a difficult problem in solving PDE's, and the above restrictions result from our approach to accomplish this.

All problems do not fit these restrictions. For example, certain hyperbolic systems of PDE's require (for stability) that the boundary conditions be imposed through the use of characteristic transformations at the boundary. We know of no way to implement the use of characteristic variables for the boundary conditions with the current structure of PDECOL. Another limitation is that PDECOL is restricted to problems in at most one space dimension. The extension to higher dimensions of the method of lines approach with either collocation or finite differences is conceptually quite simple. The basic difficulty arises from the fact that this basic approach yields a *fully implicit* discretization method. This, of course, implies that a possibly very large matrix problem must be solved. We feel certain that there are classes of PDE problems where fully implicit methods will be necessary, and for these problems extensions of these current techniques would prove fruitful. For other classes of problems, splitting or alternating direction implicit techniques may be much more efficient.

If a user's problem does not fit the prescription of the package, he has the following options: (a) find another package, (b) modify the existing package to suit his needs, or (c) modify or transform the problem to fit the package. The complexity of PDECOL and the limit of space precludes presenting details on how to modify PDECOL to accommodate problems which may not satisfy the above restrictions.

Third, the methods implemented in a program may not be adequate for a particular problem. Collocation methods (as implemented in PDECOL) will not work for all problems. For example, some PDE problems are quite sensitive to a particular conservation law being satisfied by the discrete approximation (conservation of mass, energy, particles, etc.). Collocation techniques are inherently nonconservative and difficulties can be expected if such sensitivities exist for a problem.

Summarizing, our experience in using general purpose software such as PDECOL over several years has led us to believe that general purpose software is somewhat more difficult to use because of its generality. Therefore, the probability for errors and misuse is greater and so its primary advantage (its general nature) also becomes a disadvantage. We have found that users have difficulties primarily for three reasons: (1) their problem is not well posed, (2) the methods in the general package will not work for their problem, and (3) the user has made programming errors or misused the package. Perhaps 99 percent of all difficulties fall into the third category.

6. Structure of PDECOL

Since the methods in PDECOL are based on the method of lines, PDECOL is structured much like some of the recently developed integrators for ordinary differential equations. This structure includes a driver subroutine, a core integrator which advances the time by a single time step, and miscellaneous routines which assist in setting up and solving nonlinear and linear equations.

PDECOL is somewhat more complicated than the typical ODE integrator since: (a) PDE's are being solved, and (b) a finite element method requiring the evaluation of the piecewise polynomials (4.1) is being used. The interconnections and lines of communication in PDECOL are quite complicated and difficult to understand when viewed as a whole. However, when the driver program and the core integrator are isolated, the picture becomes more clear. Figures 1 and 2 illustrate all of the package and user routines and their interconnections. Some of the routines are duplicated in each figure. A brief summary of the basic functions of all of the package subroutines is contained in the Algorithm. We will now describe the principal routines in PDECOL.

Main Components. PDECOL. This routine is a driver for the entire package. It serves the following purposes. The locations and lengths of the storage arrays

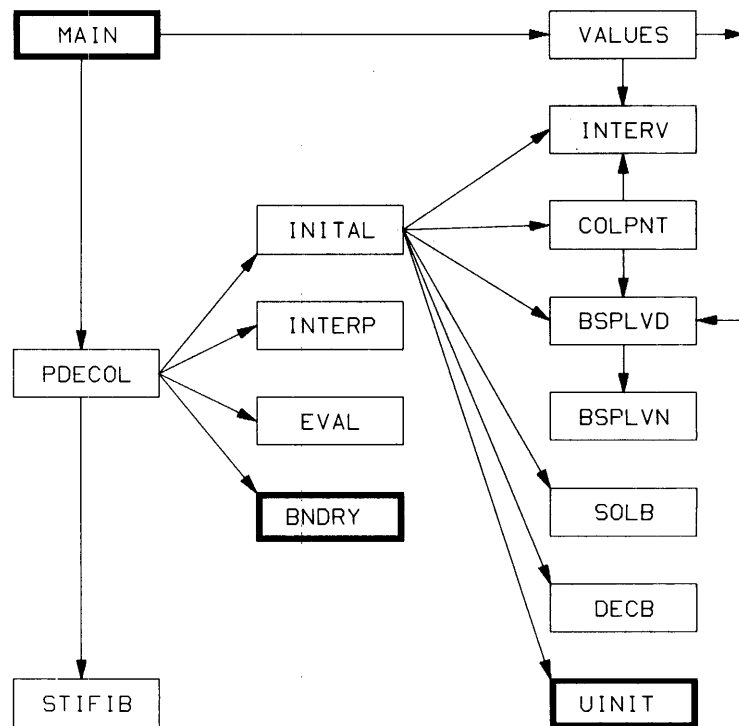


Fig. 1. Lines of communication between user routines and PDECOL package routines: Part 1, user routines are indicated by heavy borders

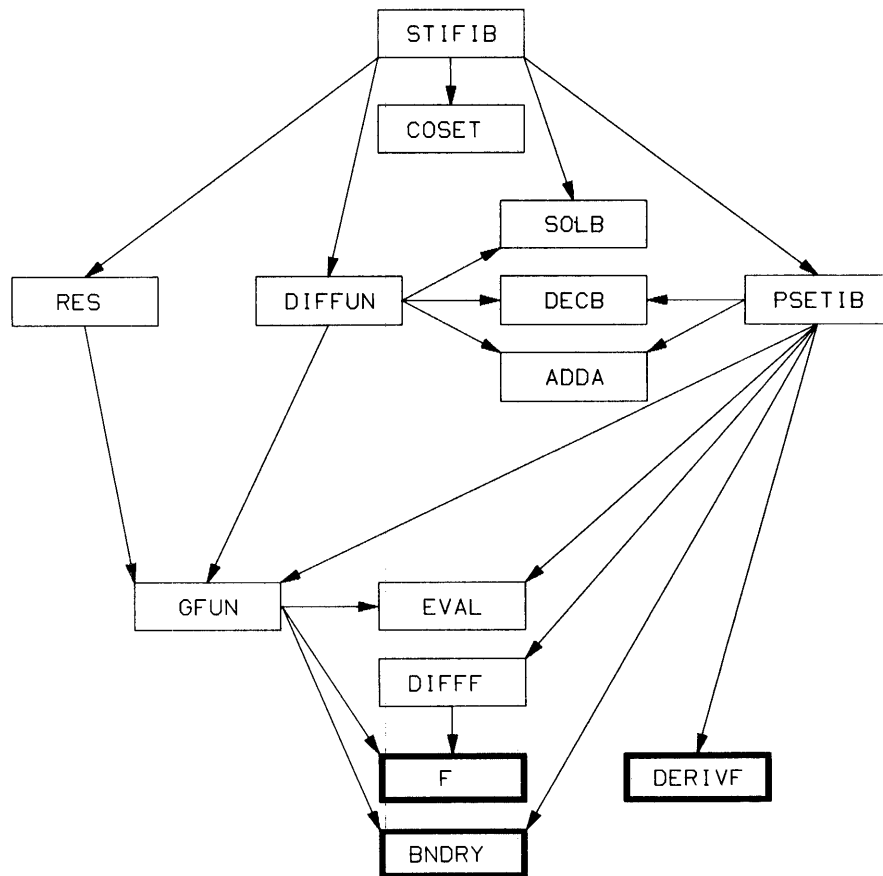


Fig. 2. Lines of communication between user routines and PDECOL package routines: Part 2, user routines are indicated by heavy borders

required by the package are allocated and defined in PDECOL. On the first call to the package, this routine checks the user's input for legality and performs initialization tasks such as setting problem parameters and calling the package routine INITAL. After the initialization is complete, PDECOL makes repeated calls to the core integrator routine STIFIB in order to advance the time variable a single step at a time until the user's desired output time is reached. At this time (since output times are not usually hit exactly) the package routine INTERP is called to interpolate the computed basis function coefficients to the desired output time. PDECOL also checks for possible error returns from STIFIB and either attempts to recover or provides the user with appropriate flags and/or messages.

VALUES. This is the routine the user must call in order to obtain the values of his computed approximate solution and its derivatives when his desired output time TOUT has been reached. VALUES performs the same basic function as EVAL (see below) except that it can evaluate the piecewise polynomial solution at *any* given spatial point or points—rather than just at collocation points. Since the computed solution is actually a continuous function, the user may desire solution values at arbitrary spatial points in the domain of interest. VALUES is structured much like EVAL. The basis function values are generated by BSPLVD as needed, and the user has the option of obtaining values at any number of points with a single call to VALUES.

INITAL. This routine performs functions which are very important to the entire package. Its main purpose is to determine the initial condition data for the package. This initial data consists of the basis function coefficients which are chosen so that the initial piecewise polynomial ((4.1) at $t = t_0$) interpolates the user's initial condition functions (2.4) at the collocation points. This requires that the collocation points and all of the basis function data be generated. INITAL generates (or causes to be generated) and stores the breakpoint sequence, the collocation points (see COLPNT), and the B-spline basis function and derivative values at the collocation points [3]. The basis function values and derivatives at the collocation points are calculated only once and saved so as to make the package efficient (see EVAL below).

COLPNT. The collocation points which are required by the collocation discretization procedure are calculated in this routine. In general we choose the collocation points to be the points in the domain $[x_L, x_R]$ at which the B-spline basis functions attain their unique maxima [1]. For the special piecewise polynomial spaces of order $KORD$ with $NCC = 2$, a second more desirable choice is available (normal default choice). In this case the collocation points are chosen to be the unique $KORD - NCC$ Gauss-Legendre quadrature points in each subinterval $[XBKPT(i), XBKPT(i + 1)]$, for $i = 1, 2, \dots, NINT$, plus the two endpoints x_L and x_R . This choice of points for these piecewise polynomial spaces will produce more accurate approximate solutions for certain classes of problems [4].

BSPLVD, BSPLVN, and INTERV. These subprograms are the basic routines which are used to generate the B-spline basis function values and their derivative values at any desired points. They were developed and fully documented by deBoor [3] and are essentially unchanged from his original implementation. They provide for the reliable evaluation of arbitrary B-spline basis function values for polynomial orders up to 20.

STIFIB. This routine is the core integrator for the package. Its basic task is to advance the time value by taking a single time step. The routines STIFIB, COSET, DIFFUN, RES, ADDA, PSETIB, INTERP, DECB, and SOLB actually constitute a slightly modified version of an ODE package called GEARIB which has been developed by Hindmarsh [6]. The package GEARIB is designed to solve initial-value ODE problems of the form

$$A \frac{dy}{dt} = \mathbf{g}(t, \mathbf{y}), \quad (6.1)$$

where $A(t, \mathbf{y})$ is a banded matrix, \mathbf{y} is a vector, and \mathbf{g} is a vector valued function. The semidiscrete equations (4.5) obtained by the collocation procedure above are particular examples of such equations. The time integration algorithms that are implemented in STIFIB are quite complex and we will discuss only the more important features. The methods used to advance the time are based on multi-point formulas of the form

$$\mathbf{y}_n = \sum_{j=1}^{k_1} \alpha_j \mathbf{y}_{n-j} + \Delta t \sum_{j=0}^{k_2} \beta_j \mathbf{y}'_{n-j}, \quad (6.2)$$

where \mathbf{y}_k is an approximation to the exact value at $t = t_k$, \mathbf{y}' denotes $d\mathbf{y}/dt$, $\Delta t = t_n - t_{n-1}$, and α_j and β_j are the method coefficients. Since $\beta_0 \neq 0$ for the methods used, the time integration methods (6.2) are *implicit* and so a nonlinear system of equations must in general be solved for every time step taken.

Proper choices of the coefficients in (6.2) can produce the standard Crank-Nicolson and backward difference methods so these techniques are generalizations of the classical, often used time integration methods.

There are two basic types of time integration formulas used in STIFIB. The first type of formula (Adams' methods) of order q , $1 \leq q \leq 12$, is obtained from (6.2) by setting $k_1 = 1$, $k_2 = q - 1$, and $\alpha_1 = 1$. The second type of formula (backward differentiation methods) of order q , $1 \leq q \leq 5$, is obtained by setting $k_1 = q$ and $k_2 = 0$. The remaining unspecified coefficients may be found in the book by Gear [5]. In general we recommend the use of the second type of formulas because of their better stability properties.

The time integration techniques in PDECOL feature automatic time step size and time integration formula order selection so as to efficiently solve the problem and yet achieve a user specified time integration error level.

For more specific details of the techniques and procedures implemented in STIFIB, we refer the reader to the comments in the program and the GEARIB report [6].

GFUN. The primary function of the routine GFUN is to properly evaluate the function $\mathbf{g}(t, \mathbf{y})$ of (6.1) when provided with input values of t and \mathbf{y} . In the package the vector function $\mathbf{g}(t, \mathbf{y})$ actually consists of values at the collocation points of the right-hand-side function \mathbf{f} , which defines the original PDE problem (2.1). Since the quantities actually being computed by the package (and the arguments input to GFUN) are basis function coefficients, GFUN must use these coefficients and the user defined subprograms F and BNDRY to generate the proper output values of $\mathbf{g}(t, \mathbf{y})$. The basis function coefficients are converted to piecewise polynomial values by EVAL (see below) and these values are then used to call the user's routines F and BNDRY. This conversion and evaluation process is performed at each collocation point in succession with BNDRY being called only for the boundary points. The first and last block rows of the A matrix (6.1) are also updated by GFUN to properly account for the boundary conditions (2.1).

GFUN is the routine which actually implements the collocation discretization process. It serves as the basic interface (via the method of lines) between the users' routines F and BNDRY and an ODE integrator (in this case a modified GEARIB). It serves the same interface role as the routine PDEONE [12] in a similar finite difference—method of lines implementation.

EVAL. The purpose of EVAL is to evaluate the piecewise polynomial (4.1) and its first two derivatives at a given collocation point, when the current basis function coefficients are provided as input. Since at any given point, at most $KORD$ basis functions are nonzero, this computation involves computing three vector inner products where the vectors are of length $KORD$. This computation is the "innermost loop" in the package, since EVAL is called just prior to each call to the user's routine F (or BNDRY). Originally, the basis function values were not stored and were repeatedly computed (by BSPLVD) as needed. However, the expense of this repeated computation was so great that overall run times

were decreased by factors of from five to ten when it was decided to store all of the basis function (and derivative) values at the collocation points (see INITAL). This gain in speed is, of course, at the expense of additional storage.

PSETIB. This routine is required by STIFIB to compute and process the matrix $A - h\beta_0 (\partial \mathbf{g} / \partial \mathbf{y})$ where A is the banded matrix in (6.1) and $\partial \mathbf{g} / \partial \mathbf{y}$ is the banded Jacobian matrix of the function $\mathbf{g}(t, \mathbf{y})$ in (6.1). If the values of the partial derivatives of the user's function \mathbf{f} in (2.1) with respect to \mathbf{u} , \mathbf{u}_x , and \mathbf{u}_{xx} are known, then the Jacobian matrix $\partial \mathbf{g} / \partial \mathbf{y}$ can easily be computed using the chain rule. Once PSETIB has formed the desired matrix, it then LU decomposes the matrix for use in STIFIB by calling the package routine DECB. There are two options available in PSETIB for generating the needed partial derivatives and the user makes the selection. One option requires that the user construct a routine DERIVF to provide the partial derivative information (see the Algorithm). In the second option approximate partial derivatives are generated internally in the package routine DIFFF by finite difference quotients obtained by calling the user's routine F. The DERIVF option is only slightly more efficient (at most 10 percent faster). However, it does have the advantage that more accurate partial derivative values will be obtained. Inaccurate Jacobian matrix values do not directly affect the accuracy of the computed approximate solution, but can cause longer running times by forcing the package to take smaller time steps and to generate Jacobian matrices more frequently. For most problems we have found the second option totally adequate and recommend its use.

7. Example Problems

To illustrate the use of PDECOL, we present a main program and the user written subroutines required to solve the following problem on the interval $[0, 1]$:

$$\begin{aligned}\frac{\partial u}{\partial t} &= v^2 \frac{\partial^2 u}{\partial x^2} + 2v \frac{\partial v}{\partial x} \frac{\partial u}{\partial x} - uv - u^2 + 10, \\ \frac{\partial v}{\partial t} &= u^2 \frac{\partial^2 v}{\partial x^2} + 2u \frac{\partial u}{\partial x} \frac{\partial v}{\partial x} + \frac{\partial^2 u}{\partial x^2} + uv - v^2,\end{aligned}$$

with boundary conditions

$$\begin{aligned}u &= \frac{1}{2}, & v &= \pi & \text{at } x = 0, \\ \frac{\partial u}{\partial x} + \sin(uv) &= \frac{1}{2}, & \frac{\partial v}{\partial x} - \cos(uv) &= 1 & \text{at } x = 1\end{aligned}$$

and initial conditions

$$u = \frac{1}{2}(x + 1), \quad v = \pi \quad \text{at } t = 0.$$

Note that the initial conditions are consistent with the boundary conditions as required by PDECOL.

Figure 3 shows a main program and the user written subroutines F, BNDRY, UINIT, and DERIVF for this problem. Subroutine DERIVF is an optional routine which needs to be provided only if $MF = 11$ or 21 (see internal computer documentation for additional details). We present it here for completeness.

In Figure 4 we present some numerical results for this problem. These results were obtained using the piecewise polynomial space defined by $KORD = 4$, $NCC = 2$, and $NINT = 30$ with equally spaced breakpoints. The reader is referred to the internal program documentation in PDECOL for the definition of the quantities that appear in the figure. Since we do not know the exact solution to this problem, we are not able to determine the error in the calculated solution.

We next consider a relatively simple PDE problem where we know the exact solution. The purpose of this problem is to show the use of higher order methods


```

PROGRAM TEST(UGOUT,TAPE3=UGOUT)
COMMON /ENDPT/ XLEFT
COMMON /GEAR0/ DTUSED,NQ,NSTEPS,NF,NJ
DIMENSION U(2,31),XBKPT(31),SCTCH(10),WORK(5000),IWORK(500)
C
C INITIALIZE PARAMETERS AND PDECOL CALLING ARGUMENTS
C
      NPDE = 2
      NINT = 30
      NPTS = NINT + 1
      KORD = 4
      NCC = 2
      T0 = 0.0
      TOUT = 1.E-3
      DT = 1.E-7
      EPS = 1.E-4
      MF = 21
      INDEX = 1
      IWORK(1) = 5000
      IWORK(2) = 500
      DX = 1.0 / FLOAT(NPTS-1)
      DO 10 I=1,NPTS
          XBKPT(I) = FLOAT(I-1) * DX
10 CONTINUE
      XLEFT = XBKPT(1)
C
C CALL THE PACKAGE TO INTEGRATE TO TIME T = TOUT
C
      20 CALL PDECOL(T0,TOUT,DT,XBKPT,EPS,NINT,KORD,NCC,NPDE,MF,INDEX
          *          WORK,IWORK)
C
C CHECK FOR EXECUTION ERRORS
C
      IF ( INDEX .NE. 0 ) GO TO 70
C
C OUTPUT PERFORMANCE DATA AND COMPUTED SOLUTION VALUES
C
      WRITE(3,30) TOUT,DTUSED,NSTEPS
30 FORMAT(/10X,3HT= ,E10.3,7H DT= ,E10.3,15H TOTAL STEPS=,I5)
      CALL VALUES(XBKPT,U,SCTCH,NPDE,NPTS,NPTS,0,WORK)
      DO 60 K=1,NPDE
          WRITE(3,40) K
40 FORMAT(/10X,13HPDE COMPONENT,I2/)
          WRITE(3,50) (U(K,I),I=1,NPTS)
50 FORMAT(10X,5E12.4)
60 CONTINUE
C
C SET NEW OUTPUT TIME AND CONTINUE THE INTEGRATION IF TOUT .LT. 11.0
C OTHERWISE, TERMINATE THE PROBLEM
C
      TOUT = TOUT * 10.0
      IF ( TOUT .LT. 11.0 ) GO TO 20
70 WRITE(3,80) INDEX
80 FORMAT(10X,7HINDEX= ,I3)
      STOP
      END

SUBROUTINE F(T,X,U,UX,UXX,FVAL,NPDE)
DIMENSION U(NPDE), US(NPDE), UXX(NPDE), FVAL(NPDE)
      FVAL(1) = U(2)*U(2)*UXX(1) - U(1)*U(2) - U(1)**2 + 10.0
      *          + 2.0*U(2)*UX(2)*UX(1)
      FVAL(2) = U(1)*U(1)*UXX(2) + U(1)*U(2) - U(2)**2
      *          + UXX(1) + 2.0*U(1)*UX(1)*UX(2)
      RETURN
      END

```

Fig. 3. Main program and user written subroutines for the first example problem

```

SUBROUTINE BNDRY(T,X,U,UX,DBDU,DBDUX,DZDT,NPDE)
DIMENSION U(NPDE), UX(NPDE), DZDT(NPDE)
DIMENSION DBDU(NPDE,NPDE), DBDUX(NPDE,NPDE)
COMMON /ENDPT/ XLEFT
IF ( X .NE. XLEFT ) GO TO 10
  DBDU(1,1) = 1.0
  DBDU(1,2) = 0.0
  DBDU(2,1) = 0.0
  DBDU(2,2) = 1.0
  DBDUX(1,1) = 0.0
  DBDUX(1,2) = 0.0
  DBDUX(2,1) = 0.0
  DBDUX(2,2) = 0.0
  DZDT(1) = 0.0
  DZDT(2) = 0.0
RETURN
10 DBDU(1,1) = U(2) * COS( U(1) * U(2) )
  DBDU(1,2) = U(1) * COS( U(1) * U(2) )
  DBDU(2,1) = U(2) * SIN( U(1) * U(2) )
  DBDU(2,2) = U(1) * SIN( U(1) * U(2) )
  DBDUX(1,1) = 1.0
  DBDUX(1,2) = 0.0
  DBDUX(2,1) = 0.0
  DBDUX(2,2) = 1.0
  DZDT(1) = 0.0
  DZDT(2) = 0.0
RETURN
END

SUBROUTINE UINIT(X,U,NPDE)
DIMENSION U(NPDE)
C
C SET INITIAL CONDITIONS. NOTE THAT PI = 4.0*ATAN(1.0)
C
  U(1) = 0.5 * ( X = 1.0 )
  U(2) = 4.0 * ATAN( 1.0 )

RETURN
END

SUBROUTINE DERIVF(T,X,U,UX,UXX,DFDU,DFDUX,DFDUXX,NPDE)
DIMENSION U(NPDE), UX(NPDE), UXX(NPDE)
DIMENSION DFDU(NPDE,NPDE), DFDUX(NPDE,NPDE), DFDUXX(NPDE,NPDE)
  DFDU(1,1) = -U(2) - 2.0*U(1)
  DFDU(1,2) = 2.0*U(2)*UXX(1) - U(1) + 2.0*UX(2)*UX(1)
  DFDU(2,1) = 2.0*U(1)*UXX(2) + U(2) + 2.0*UX(1)*UX(2)
  DFDU(2,2) = U(1) - 2.0*U(2)
C
  DFDUX(1,1) = 2.0*U(2)*UX(2)
  DFDUX(1,2) = 2.0*U(2)*UX(1)
  DFDUX(2,1) = 2.0*U(1)*UX(2)
  DFDUX(2,2) = 2.0*U(1)*UX(1)
C
  DFDUXX(1,1) = U(2)*U(2)
  DFDUXX(1,2) = 0.0
  DFDUXX(2,1) = 1.0
  DFDUXX(2,2) = U(1)*U(1)
RETURN
END

```

Fig. 3. (Continued)

and the execution times on a CDC-7600 computer. The problem is of a simple diffusion type and is described by

$$\frac{\partial u}{\partial t} = \frac{\partial^2 u}{\partial x^2} + \pi^2 \sin \pi x,$$

$$u(0, x) = 1, \quad u(t, 0) = u(t, 1) = 1$$

for $0 < x < 1$ and $t > 0$. The exact solution for this problem is easily seen to be

$$u(t, x) = 1 + (\sin \pi x)(1 - e^{-\pi^2 t}).$$

Table I shows the results obtained for this problem at $t = 0.1$ with $NCC = 2$ for several values of $KORD$ and $NINT$. Equally spaced breakpoints and $MF = 21$ were used in all cases. The values of EPS , the time integration error tolerance, were determined by *trial and error* to be as large as possible without affecting the spatial errors. The initial DT was chosen to be equal to $EPS/2$. $NSTEPS$ is the number of time integration steps. The time, $TIME$, is the computer time in seconds (CDC-7600) required to produce the solution. The error, $ERROR$, is the maximum absolute difference between the exact solution and the numerical solution. The calculated order of convergence is shown in the last column. The numbers in parentheses are the theoretical values.

The preceding two examples illustrate the use of PDECOL and demonstrate the qualitative behavior of the higher order collocation methods that are implemented in the package. To consider all combinations of PDE methods that are built into PDECOL is virtually impossible since there are potentially over 3000 different methods in the package. We refer the reader to several other papers [9, 10] for additional results on the use, testing, and evaluation of PDECOL to solve partial differential equations and two-point boundary value problems. Some comparisons of PDECOL with low order finite difference methods are contained

$T=$	$1.000E-03$	$DT=$	$3.400E-04$	$TOTAL$	$STEPS=$	11
PDE COMPONENT 1						
$5.0000E-01$	$5.1915E-01$	$5.3759E-01$	$5.5547E-01$	$5.7294E-01$	$5.9012E-01$	$6.0709E-01$
$6.7402E-01$	$6.9065E-01$	$7.0726E-01$	$7.2386E-01$	$7.4046E-01$	$7.5706E-01$	$7.7365E-01$
$8.4002E-01$	$8.5661E-01$	$8.7321E-01$	$8.8982E-01$	$9.0645E-01$	$9.2309E-01$	$9.3977E-01$
$1.0074E+00$						
PDE COMPONENT 2						
$3.1416E+00$	$3.1335E+00$	$3.1331E+00$	$3.1332E+00$	$3.1334E+00$	$3.1335E+00$	$3.1336E+00$
$3.1338E+00$	$3.1339E+00$	$3.1340E+00$	$3.1340E+00$	$3.1341E+00$	$3.1341E+00$	$3.1342E+00$
$3.1344E+00$	$3.1344E+00$	$3.1345E+00$	$3.1345E+00$	$3.1346E+00$	$3.1346E+00$	$3.1347E+00$
$3.1347E+00$	$3.1347E+00$	$3.1348E+00$	$3.1349E+00$	$3.1349E+00$	$3.1350E+00$	
$3.1350E+00$						
$T=$	$1.000E-02$	$DT=$	$2.052E-03$	$TOTAL$	$STEPS=$	21
PDE COMPONENT 1						
$5.0000E-01$	$5.2478E-01$	$5.4920E-01$	$5.7300E-01$	$5.9609E-01$	$6.1848E-01$	$6.4020E-01$
$7.2165E-01$	$7.4097E-01$	$7.5999E-01$	$7.7877E-01$	$7.9738E-01$	$8.1588E-01$	$8.3433E-01$
$9.0902E-01$	$9.2829E-01$	$9.4797E-01$	$9.6817E-01$	$9.8898E-01$	$1.0105E+00$	$1.0330E+00$
$1.1345E+00$						

Fig. 4. Numerical results for the first example problem (continued on p. 342)

```

PDECOMPONENT 2
  3.1416E+00  3.0943E+00  3.0722E+00  3.0634E+00  3.0606E+00
  3.0604E+00  3.0611E+00  3.0620E+00  3.0630E+00  3.0640E+00
  3.0649E+00  3.0658E+00  3.0667E+00  3.0675E+00  3.0684E+00
  3.0692E+00  3.0700E+00  3.0708E+00  3.0716E+00  3.0724E+00
  3.0733E+00  3.0742E+00  3.0752E+00  3.0763E+00  3.0773E+00
  3.0785E+00  3.0798E+00  3.0811E+00  3.0825E+00  3.0842E+00
  3.0861E+00

```

T= 1.000E-01 DT= 7.857E-03 TOTAL STEPS= 52

```

PDE COMPONENT 1
  5.0000E-01  5.4709E-01  5.9690E-01  6.4852E-01  7.0117E-01
  7.5419E-01  8.0704E-01  8.5929E-01  9.1062E-01  9.6079E-01
  1.0096E+00  1.0570E+00  1.1028E+00  1.1471E+00  1.1897E+00
  1.2308E+00  1.2702E+00  1.3081E+00  1.3444E+00  1.3793E+00
  1.4127E+00  1.4447E+00  1.4753E+00  1.5047E+00  1.5328E+00
  1.5598E+00  1.5855E+00  1.6102E+00  1.6339E+00  1.6566E+00
  1.6783E+00

```

```

PDE COMPONENT 2
  3.1416E+00  3.0064E+00  2.9094E+00  2.8409E+00  2.7943E+00
  2.7643E+00  2.7475E+00  2.7409E+00  2.7428E+00  2.7514E+00
  2.7657E+00  2.7847E+00  2.8079E+00  2.8347E+00  2.8647E+00
  2.8977E+00  2.9332E+00  2.9713E+00  3.0117E+00  3.0543E+00
  3.0991E+00  3.1459E+00  3.1948E+00  3.2457E+00  3.2986E+00
  3.3535E+00  3.4104E+00  3.4693E+00  3.5303E+00  3.5933E+00
  3.6584E+00

```

T= 1.000E+00 DT= 3.566E-01 TOTAL STEPS= 76

```

PDE COMPONENT 1
  5.0000E-01  5.4699E-01  5.9679E-01  6.4846E-01  7.0117E-01
  7.5420E-01  8.0695E-01  8.5895E-01  9.0965E-01  9.5940E-01
  1.0074E+00  1.0538E+00  1.0985E+00  1.1414E+00  1.1826E+00
  1.2221E+00  1.2599E+00  1.2962E+00  1.3308E+00  1.3640E+00
  1.3957E+00  1.4261E+00  1.4551E+00  1.4830E+00  1.5096E+00
  1.5351E+00  1.5596E+00  1.5830E+00  1.6055E+00  1.6271E+00
  1.6479E+00

```

```

PDE COMPONENT 2
  3.1416E+00  3.0030E+00  2.9035E+00  2.8341E+00  2.7879E+00
  2.7598E+00  2.7460E+00  2.7435E+00  2.7500E+00  2.7638E+00
  2.7837E+00  2.8085E+00  2.8374E+00  2.8699E+00  2.9054E+00
  2.9436E+00  2.9842E+00  3.0270E+00  3.0718E+00  3.1184E+00
  3.1669E+00  3.2170E+00  3.2688E+00  3.3223E+00  3.3774E+00
  3.4341E+00  3.4925E+00  3.5525E+00  3.6143E+00  3.6778E+00
  3.7432E+00

```

T= 1.000E+01 DT= 3.566E+00 TOTAL STEPS= 81

```

PDE COMPONENT 1
  5.0000E-01  5.4699E-01  5.9679E-01  6.4846E-01  7.0117E-01
  7.5420E-01  8.0695E-01  8.5895E-01  9.0965E-01  9.5940E-01
  1.0074E+00  1.0538E+00  1.0985E+00  1.1414E+00  1.1826E+00
  1.2221E+00  1.2599E+00  1.2962E+00  1.3308E+00  1.3640E+00
  1.3957E+00  1.4261E+00  1.4551E+00  1.4830E+00  1.5096E+00
  1.5351E+00  1.5596E+00  1.5830E+00  1.6055E+00  1.6271E+00
  1.6479E+00

```

```

PDE COMPONENT 2
  3.1416E+00  3.0030E+00  2.9035E+00  2.8341E+00  2.7879E+00
  2.7598E+00  2.7460E+00  2.7434E+00  2.7500E+00  2.7638E+00
  2.7837E+00  2.8085E+00  2.8374E+00  2.8699E+00  2.9054E+00
  2.9436E+00  2.9842E+00  3.0270E+00  3.0718E+00  3.1184E+00
  3.1669E+00  3.2170E+00  3.2688E+00  3.3223E+00  3.3774E+00
  3.4341E+00  3.4925E+00  3.5525E+00  3.6143E+00  3.6778E+00
  3.7432E+00

```

INDEX= 0

Fig. 4. (Continued)

Table I. Results at $t = 0.1$ for the Second Example Problem with $NCC = 2$ for Several Values of $KORD$ and $NINT$

$KORD$	$NINT$	EPS	$NSTEPS$	$TIME$ (sec)	$ERROR$	Order of convergence
3	10	3×10^{-5}	15	3.33×10^{-2}	1.14×10^{-3}	
3	20	1×10^{-5}	19	7.56×10^{-2}	2.94×10^{-4}	1.96 (2)
4	8	1×10^{-6}	25	9.53×10^{-2}	2.19×10^{-5}	
4	16	1×10^{-7}	33	2.32×10^{-1}	1.53×10^{-6}	3.90 (4)
5	4	1×10^{-6}	25	8.44×10^{-2}	1.38×10^{-5}	
5	8	1×10^{-8}	44	2.50×10^{-1}	4.33×10^{-7}	5.00 (5)
6	2	1×10^{-6}	25	6.61×10^{-2}	1.85×10^{-5}	
6	4	1×10^{-8}	41	1.84×10^{-1}	4.13×10^{-7}	5.50 (6)

in [9]. Our basic feelings are that for *smooth* problems and *high* accuracy demands, the higher order collocation techniques will be significantly more efficient, that is, produce more accuracy per unit of run time.

Our experience in using PDECOL has led us to firmly believe that PDECOL is a versatile software package that can solve a broad class of nonlinear partial differential equations. Its higher order methods can produce extremely accurate solutions quite efficiently when compared to lower order methods. The package is quite portable.

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ALGORITHM

[Summary information and a part of the listing is printed here. To facilitate understanding of the partial listing given here, lines 6110-6180 of the complete listing have been substituted for lines 4050-4060 of the complete listing, which is available from the ACM Algorithms Distribution Service.]

- NAME(n): indicates a Fortran module with n records
NAME^D(n): indicates "NAME" contains test data
NAME^M(n): indicates "NAME" contains machine dependent data

Contents: PDECOL (924), VALUES (70), BLOCKD^M (19), INITAL (84), COLPNT (222), BSPLVD (72), BSPLVN (46), INTERV (88), STIFIB (416), GFUN (70), EVAL (28), DIFFUN (28), ADDA (45), RES (53), PSETIB (102), DIFFF (46), INTERP (31), COSET (179), DECB (87), SOLB (47)

Test package: TEST1 (56), F1 (8), BNDRY1 (28), UINITI (9), DERIVF1 (19), TEST2 (61), F2 (6), BNDRY2 (6), UNINIT2 (5), DERIVF2 (6), TRUSOL2 (5), TSIN2^D (21), TEST3 (61), F3 (10), BNDRY3 (27), UINIT3 (5), DERIVF3 (17), TRUSOL3 (6), TSIN3^D (25), TEST4 (64), F4 (5), BNDRY4 (12), UINIT4 (5), DERIVF4 (6), TRUSOL4 (11), TSIN4^D (19), TEST5 (61), F5 (5), BNDRY5 (13), UINIT5 (5), DERIVF5 (6), TRUSOL5 (5), TSIN5^D (24)

```

SUBROUTINE PDECOL(T0,TOUT,DT,XBKPT,EPS,NINT,KORD,NCC,NPDE,MF,
*          INDEX,WORK,IWORK)
C
C-----
C-----
C
C THIS IS THE MARCH 24, 1978 VERSION OF PDECOL.
C
C THIS PACKAGE WAS CONSTRUCTED SO AS TO CONFORM TO AS MANY ANSI-FORTRAN
C RULES AS WAS CONVENIENTLY POSSIBLE. THE FORTRAN USED VIOLATES ANSI
C STANDARDS IN THE TWO WAYS LISTED BELOW....
C
C 1. SUBSCRIPTS OF THE GENERAL FORM C*V1 + V2 + V3 ARE USED
C (POSSIBLY IN A PERMUTED ORDER), WHERE C IS AN INTEGER CONSTANT
C AND V1, V2, AND V3 ARE INTEGER VARIABLES.
C
C 2. ARRAY NAMES APPEAR SINGLY IN DATA STATEMENTS IN THE ROUTINES
C BSPLVN AND COSET.
C
C MACHINE DEPENDENT FEATURES.....
C
C THIS VERSION OF PDECOL WAS DESIGNED FOR USE ON CDC MACHINES WITH
C A WORD LENGTH OF 60 BITS. WE DO NOT RECOMMEND THE USE OF PDECOL WITH
C WORD LENGTHS OF LESS THAN 48 BITS. THE MOST IMPORTANT MACHINE
C AND WORD LENGTH DEPENDENT CONSTANTS ARE DEFINED IN THE BLOCK DATA
C AND IN SUBROUTINES COLPNT AND COSET. THE USER SHOULD CHECK THESE
C CAREFULLY FOR APPROPRIATENESS FOR HIS LOCAL SITUATION. THE FORTRAN
C FUNCTIONS USED BY EACH ROUTINE ARE LISTED IN THE COMMENTS TO
C FACILITATE CONVERSION TO DOUBLE PRECISION.
C
C-----
C-----
C
C PDECOL IS THE DRIVER ROUTINE FOR A SOPHISTICATED PACKAGE OF
C SUBROUTINES WHICH IS DESIGNED TO SOLVE THE GENERAL SYSTEM OF
C NPDE NONLINEAR PARTIAL DIFFERENTIAL EQUATIONS OF AT MOST SECOND
C ORDER ON THE INTERVAL (XLEFT,XRIGHT) FOR T .GT. T0 WHICH IS OF THE
C FORM....
C
C      DU/DT = F( T, X, U, UX, UXX )
C
C WHERE
C
C      U = ( U(1), U(2), ... , U(NPDE) )
C      UX = ( UX(1), UX(2), ... , UX(NPDE) )
C      UXX = ( UXX(1), UXX(2), ... , UXX(NPDE) ) .
C
C EACH U(K) IS A FUNCTION OF THE SCALAR QUANTITIES T AND X.
C UX(K) REPRESENTS THE FIRST PARTIAL DERIVATIVE OF U(K) WITH RESPECT
C TO THE VARIABLE X, UXX(K) REPRESENTS THE SECOND PARTIAL DERIVATIVE
C OF U(K) WITH RESPECT TO THE VARIABLE X, AND DU/DT IS THE VECTOR OF
C PARTIAL DERIVATIVES OF U WITH RESPECT TO THE TIME VARIABLE T.
C F REPRESENTS AN ARBITRARY VECTOR VALUED FUNCTION WHOSE NPDE
C COMPONENTS DEFINE THE RESPECTIVE PARTIAL DIFFERENTIAL EQUATIONS OF
C THE PDE SYSTEM. SEE SUBROUTINE F DESCRIPTION BELOW.
C
C BOUNDARY CONDITIONS
C

```

```

C   DEPENDING ON THE TYPE OF PDE(S),  $\phi$ , 1, OR 2 BOUNDARY CONDITIONS      600
C   ARE REQUIRED FOR EACH PDE IN THE SYSTEM. THESE ARE IMPOSED AT XLEFT    610
C   AND/OR XRIGHT AND EACH MUST BE OF THE FORM...                          620
C                                                                           630
C       B(U,UX) = Z(T)                                                    640
C                                                                           650
C   WHERE B AND Z ARE ARBITRARY VECTOR VALUED FUNCTIONS WITH              660
C   NPDE COMPONENTS AND U, UX, AND T ARE AS ABOVE. THESE BOUNDARY        670
C   CONDITIONS MUST BE CONSISTENT WITH THE INITIAL CONDITIONS WHICH ARE    680
C   DESCRIBED NEXT.                                                       690
C                                                                           700
C   INITIAL CONDITIONS                                                    710
C                                                                           720
C   EACH SOLUTION COMPONENT U(K) IS ASSUMED TO BE A KNOWN (USER          730
C   PROVIDED) FUNCTION OF X AT THE INITIAL TIME T = T0. THE              740
C   INITIAL CONDITION FUNCTIONS MUST BE CONSISTENT WITH THE BOUNDARY      750
C   CONDITIONS ABOVE, I.E. THE INITIAL CONDITION FUNCTIONS MUST          760
C   SATISFY THE BOUNDARY CONDITIONS FOR T = T0. SEE SUBROUTINE UINIT      770
C   DESCRIPTION BELOW.                                                    780
C-----                                                                    790
C                                                                           800
C   REQUIRED USER SUPPLIED SUBROUTINES                                     810
C                                                                           820
C   THE USER IS REQUIRED TO CONSTRUCT THREE SUBPROGRAMS AND A MAIN        830
C   PROGRAM WHICH DEFINE THE PDE PROBLEM WHOSE SOLUTION IS TO BE        840
C   ATTEMPTED. THE THREE SUBPROGRAMS ARE...                               850
C                                                                           860
C 1) SUBROUTINE F( T, X, U, UX, UXX, FVAL, NPDE )                          870
C   DIMENSION U(NPDE), UX(NPDE), UXX(NPDE), FVAL(NPDE)                   880
C   THIS ROUTINE DEFINES THE DESIRED PARTIAL DIFFERENTIAL                890
C   EQUATIONS TO BE SOLVED. THE PACKAGE PROVIDES VALUES OF THE          900
C   INPUT SCALARS T AND X AND INPUT ARRAYS (LENGTH NPDE) U, UX,          910
C   AND UXX, AND THE USER MUST CONSTRUCT THIS ROUTINE TO COMPUTE        920
C   THE OUTPUT ARRAY FVAL (LENGTH NPDE) WHICH CONTAINS THE              930
C   CORRESPONDING VALUES OF THE RIGHT HAND SIDES OF THE DESIRED        940
C   PARTIAL DIFFERENTIAL EQUATIONS, I.E.                                  950
C                                                                           960
C       FVAL(K) = THE VALUE OF THE RIGHT HAND SIDE OF THE K-TH PDE IN     970
C       THE PDE SYSTEM ABOVE, FOR K = 1 TO NPDE.                          980
C                                                                           990
C       THE INCOMING VALUE OF THE SCALAR QUANTITY X WILL BE A           1000
C       COLLOCATION POINT VALUE (SEE INITAL AND COLPNT) AND THE           1010
C       INCOMING VALUES IN THE ARRAYS U, UX AND UXX CORRESPOND TO THIS  1020
C       POINT X AND TIME T.                                              1030
C   RETURN                                                                1040
C   END                                                                    1050
C                                                                           1060
C 2) SUBROUTINE BNDRY( T, X, U, UX, DBDU, DBDUX, DZDT, NPDE )           1070
C   DIMENSION U(NPDE), UX(NPDE), DZDT(NPDE)                               1080
C   DIMENSION DBDU(NPDE,NPDE), DBDUX(NPDE,NPDE)                          1090
C   THIS ROUTINE IS USED TO PROVIDE THE PDE PACKAGE WITH NEEDED          1100
C   INFORMATION ABOUT THE BOUNDARY CONDITION FUNCTIONS B AND Z           1110
C   ABOVE. THE PACKAGE PROVIDES VALUES OF THE INPUT VARIABLES           1120
C   T, X, U, AND UX, AND THE USER IS TO DEFINE THE CORRESPONDING       1130
C   OUTPUT VALUES OF THE DERIVATIVES OF THE FUNCTIONS B AND Z          1140
C   WHERE...                                                              1150
C       DBDU(K,J) = PARTIAL DERIVATIVE OF THE K-TH COMPONENT OF THE      1160
C       VECTOR FUNCTION B(U,UX) ABOVE WITH RESPECT TO                    1170
C       THE J-TH VARIABLE U(J).                                          1180
C       DBDUX(K,J) = PARTIAL DERIVATIVE OF THE K-TH COMPONENT OF THE     1190
C       VECTOR FUNCTION B(U,UX) ABOVE WITH RESPECT TO                    1200
C       THE J-TH VARIABLE UX(J).                                         1210
C       DZDT(K) = DERIVATIVE OF THE K-TH COMPONENT OF THE VECTOR         1220
C       FUNCTION Z(T) ABOVE WITH RESPECT TO THE                          1230
C       VARIABLE T.                                                       1240
C   NOTE... THE INCOMING VALUE OF X WILL BE EITHER XLEFT OR XRIGHT.     1250
C   IF NO BOUNDARY CONDITION IS DESIRED FOR SAY THE K-TH PDE AT         1260
C   ONE OR BOTH OF THE ENDPOINTS XLEFT OR XRIGHT, THEN DBDU(K,K)        1270
C   AND DBDUX(K,K) SHOULD BOTH BE SET TO ZERO WHEN BNDRY IS             1280
C   CALLED FOR THAT POINT. WE REFER TO THIS AS A NULL BOUNDARY          1290
C   CONDITION. THIS ROUTINE CAN BE STRUCTURED AS FOLLOWS...             1300
C   THE COMMON BLOCK /ENDPT/ IS NOT A PART OF PDECOL AND                1310
C   MUST BE SUPPLIED AND DEFINED BY THE USER.                           1320
C   COMMON /ENDPT/ XLEFT                                                  1330
C   IF( X .NE. XLEFT ) GO TO 10                                          1340
C   HERE DEFINE AND SET PROPER VALUES FOR DBDU(K,J), DBDUX(K,J),       1350

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C   NINT = ENOUGH SO THAT ANY FINE STRUCTURE OF THE PROBLEM MAY BE      2910
C   RESOLVED.                                                            2920
C   KORD = 4                                                              2930
C   NCC = 2                                                               2940
C   MF = 22                                                               2950
C   INDEX = 1 (ON FIRST CALL ONLY, THEN 0 THEREAFTER).                 2960
C                                                                           2970
C                                                                           2980
C THE INPUT PARAMETERS ARE..                                             2990
C   T0 = THE INITIAL VALUE OF T, THE INDEPENDENT VARIABLE              3000
C   (USED ONLY ON FIRST CALL).                                          3010
C   TOUT = THE VALUE OF T AT WHICH OUTPUT IS DESIRED NEXT. SINCE      3020
C   THE PACKAGE CHOOSES ITS OWN TIME STEP SIZES, THE                    3030
C   INTEGRATION WILL NORMALLY GO SLIGHTLY BEYOND TOUT                   3040
C   AND THE PACKAGE WILL INTERPOLATE TO T = TOUT.                       3050
C   DT = THE INITIAL STEP SIZE IN T, IF INDEX = 1, OR, THE             3060
C   MAXIMUM STEP SIZE ALLOWED (MUST BE .GT. 0), IF INDEX = 3.          3070
C   USED FOR INPUT ONLY WHEN INDEX = 1 OR 3. SEE BELOW.                 3080
C   XBKPT = THE ARRAY OF PIECEWISE POLYNOMIAL BREAKPOINTS.              3090
C   THE NINT+1 VALUES MUST BE STRICTLY INCREASING WITH                 3100
C   XBKPT(1) = XLEFT AND XBKPT(NINT+1) = XRIGHT (USED ONLY             3110
C   ON FIRST CALL).                                                      3120
C   EPS = THE RELATIVE TIME ERROR BOUND (USED ONLY ON THE              3130
C   FIRST CALL, UNLESS INDEX = 4). SINGLE STEP ERROR                    3140
C   ESTIMATES DIVIDED BY CMAX(I) WILL BE KEPT LESS THAN                 3150
C   EPS IN ROOT-MEAN-SQUARE NORM. THE VECTOR CMAX OF WEIGHTS           3160
C   IS COMPUTED IN PDECOL. INITIALLY CMAX(I) IS SET TO                   3170
C   ABS(C(I)), WITH A DEFAULT VALUE OF 1 IF ABS(C(I)) .LT. 1.           3180
C   THEREAFTER, CMAX(I) IS THE LARGEST VALUE                             3190
C   OF ABS(C(I)) SEEN SO FAR, OR THE INITIAL CMAX(I) IF                 3200
C   THAT IS LARGER. TO ALTER EITHER OF THESE, CHANGE THE                 3210
C   APPROPRIATE STATEMENTS IN THE DO-LOOPS ENDING AT                    3220
C   STATEMENTS 50 AND 130 BELOW. THE USER SHOULD EXERCISE              3230
C   SOME DISCRETION IN CHOOSING EPS. IN GENERAL, THE                     3240
C   OVERALL RUNNING TIME FOR A PROBLEM WILL BE GREATER IF               3250
C   EPS IS CHOSEN SMALLER. THERE IS USUALLY LITTLE REASON TO           3260
C   CHOOSE EPS MUCH SMALLER THAN THE ERRORS WHICH ARE BEING            3270
C   INTRODUCED BY THE USERS CHOICE OF THE POLYNOMIAL SPACE.             3280
C   SEE RELATED COMMENTS CONCERNING CMAX BELOW STATEMENT 40.            3290
C   NINT = THE NUMBER OF SUBINTERVALS INTO WHICH THE SPATIAL DOMAIN      3300
C   (XLEFT,XRIGHT) IS TO BE DIVIDED (MUST BE .GE. 1)                    3310
C   (USED ONLY ON FIRST CALL).                                          3320
C   KORD = THE ORDER OF THE PIECEWISE POLYNOMIAL SPACE TO BE USED.       3330
C   ITS VALUE MUST BE GREATER THAN 2 AND LESS THAN 21. FOR              3340
C   FIRST ATTEMPTS WE RECOMMEND KORD = 4. IF THE SOLUTION                3350
C   IS SMOOTH AND MUCH ACCURACY IS DESIRED, HIGHER VALUES              3360
C   MAY PROVE TO BE MORE EFFICIENT. WE HAVE SELDOM USED                 3370
C   VALUES OF KORD IN EXCESS OF 8 OR 9, THOUGH THEY ARE               3380
C   AVAILABLE FOR USE IN PDECOL (USED ONLY ON FIRST CALL).              3390
C   NCC = THE NUMBER OF CONTINUITY CONDITIONS TO BE IMPOSED ON THE       3400
C   APPROXIMATE SOLUTION AT THE BREAKPOINTS IN XBKPT.                   3410
C   NCC MUST BE GREATER THAN 1 AND LESS THAN KORD. WE                   3420
C   RECOMMEND THE USE OF NCC = 2                                         3430
C   SINCE THEORY PREDICTS THAT DRAMATICALLY MORE                        3440
C   ACCURATE RESULTS CAN OFTEN BE OBTAINED USING THIS CHOICE            3450
C   (USED ONLY ON FIRST CALL).                                          3460
C   NPDE = THE NUMBER OF PARTIAL DIFFERENTIAL EQUATIONS IN THE SYSTEM    3470
C   TO BE SOLVED (USED ONLY ON FIRST CALL).                              3480
C   MF = THE METHOD FLAG (USED ONLY ON FIRST CALL, UNLESS                 3490
C   INDEX = 4). ALLOWED VALUES ARE 11, 12, 21, 22.                     3500
C   FOR FIRST ATTEMPTS WE RECOMMEND THE USE OF MF = 22.                 3510
C   MF HAS TWO DECIMAL DIGITS, METH AND MITER                            3520
C   (MF = 10*METH + MITER).                                              3530
C   METH IS THE BASIC METHOD INDICATOR..                                  3540
C   METH = 1 MEANS THE ADAMS METHODS (GENERALIZATIONS OF                 3550
C   CRANK-NICOLSON).                                                    3560
C   METH = 2 MEANS THE BACKWARD DIFFERENTIATION                         3570
C   FORMULAS (BDF), OR STIFF METHODS OF GEAR.                           3580
C   MITER IS THE ITERATION METHOD INDICATOR                               3590
C   AND DETERMINES HOW THE JACOBIAN MATRIX IS                            3600
C   TO BE COMPUTED..                                                    3610
C   MITER = 1 MEANS CHORD METHOD WITH ANALYTIC JACOBIAN.                 3620
C   FOR THIS USER SUPPLIES SUBROUTINE DERIVF.                            3630
C   SEE DESCRIPTION ABOVE.                                              3640
C   MITER = 2 MEANS CHORD METHOD WITH JACOBIAN CALCULATED                 3650
C   INTERNALLY BY FINITE DIFFERENCES. SEE                                3660

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C          SUBROUTINES PSETIB AND DIFFF.          3670
C INDEX = INTEGER USED ON INPUT TO INDICATE TYPE OF CALL, 3680
C          WITH THE FOLLOWING VALUES AND MEANINGS.. 3690
C          1 THIS IS THE FIRST CALL FOR THIS PROBLEM. 3700
C          0 THIS IS NOT THE FIRST CALL FOR THIS PROBLEM, 3710
C            AND INTEGRATION IS TO CONTINUE. 3720
C          2 SAME AS 0 EXCEPT THAT TOUT IS TO BE HIT 3730
C            EXACTLY (NO INTERPOLATION IS DONE). SEE NOTE 3740
C            BELOW. ASSUMES TOUT .GE. THE CURRENT T. 3750
C            IF TOUT IS .LT. THE CURRENT TIME, THEN TOUT IS 3760
C            RESET TO THE CURRENT TIME AND CONTROL IS 3770
C            RETURNED TO THE USER. A CALL TO VALUES WILL 3780
C            PRODUCE ANSWERS FOR THE NEW VALUE OF TOUT. 3790
C          3 SAME AS 0 EXCEPT CONTROL RETURNS TO CALLING 3800
C            PROGRAM AFTER ONE STEP. TOUT IS IGNORED AND 3810
C            DT MUST BE SET .GT. 0 TO A MAXIMUM ALLOWED 3820
C            DT VALUE. SEE ABOVE. 3830
C          4 THIS IS NOT THE FIRST CALL FOR THE PROBLEM, 3840
C            AND THE USER HAS RESET EPS AND/OR MF. 3850
C            SINCE THE NORMAL OUTPUT VALUE OF INDEX IS 0, 3860
C            IT NEED NOT BE RESET FOR NORMAL CONTINUATION. 3870
C 3880
C NOTE.. THE PACKAGE MUST HAVE TAKEN AT LEAST ONE SUCCESSFUL TIME 3890
C STEP BEFORE A CALL WITH INDEX = 2 OR 4 IS ALLOWED. 3900
C AFTER THE INITIAL CALL, IF A NORMAL RETURN OCCURRED AND A NORMAL 3910
C CONTINUATION IS DESIRED, SIMPLY RESET TOUT AND CALL AGAIN. 3920
C ALL OTHER PARAMETERS WILL BE READY FOR THE NEXT CALL. 3930
C A CHANGE OF PARAMETERS WITH INDEX = 4 CAN BE MADE AFTER 3940
C EITHER A SUCCESSFUL OR AN UNSUCCESSFUL RETURN PROVIDED AT LEAST 3950
C ONE SUCCESSFUL TIME STEP HAS BEEN MADE. 3960
C 3970
C WORK = FLOATING POINT WORKING ARRAY FOR PDECOL. WE RECOMMEND 3980
C          THAT IT BE INITIALIZED TO ZERO BEFORE THE FIRST CALL 3990
C          TO PDECOL. ITS TOTAL LENGTH MUST BE AT LEAST 4000
C 4010
C          KORD + 4*NPDE + 9*NPDE**2 + NCPTS*(3*KCRD + 2) + 4020
C          NPDE*NCPTS*(3*ML + MAXDER + 7) 4030
C 4040
C WHERE... 6110
C 6120
C          NCPTS = KORD*NINT - NCC*(NINT-1) 6130
C          ML = NPDE*(KORD+IQUAD-1) - 1 6140
C          IQUAD = 1 IF KORD = 3 AND A NULL BOUNDARY CONDITION EXISTS 6150
C          IQUAD = 0 OTHERWISE 6160
C          MAXDER = 5 UNLESS OTHERWISE SET BY THE USER INTO /OPTION/. 6170
C 6180
C IWORK = INTEGER WORKING ARRAY FOR PDECOL. THE FIRST TWO 4070
C LOCATIONS MUST BE DEFINED AS FOLLOWS... 4080
C IWORK(1) = LENGTH OF USERS ARRAY WORK 4090
C IWORK(2) = LENGTH OF USERS ARRAY IWORK 4100
C THE TOTAL LENGTH OF IWORK MUST BE AT LEAST 4110
C NCPTS*(NPDE + 1). 4120
C OUTPUT 4130
C 4140
C THE SOLUTION VALUES ARE NOT RETURNED DIRECTLY TO THE USER BY PDECOL. 4150
C THE METHODS USED IN PDECOL COMPUTE BASIS FUNCTION COEFFICIENTS, SO 4160
C THE USER (AFTER A RETURN FROM PDECOL) MUST CALL THE PACKAGE ROUTINE 4170
C VALUES TO OBTAIN HIS APPROXIMATE SOLUTION VALUES AT ANY DESIRED SPACE 4180
C POINTS X AT THE TIME T = TOUT. SEE THE COMMENTS IN SUBROUTINE VALUES 4190
C FOR DETAILS ON HOW TO PROPERLY MAKE THE CALL. 4200
C 4210
C EXECUTION ERROR MESSAGES WILL BE PRINTED BY PDECOL ON LOGICAL UNIT 4220
C LOUT WHICH IS THE ONLY VARIABLE IN THE COMMON BLOCK /IOUNIT/. A 4230
C DEFAULT OF LOUT = 3 IS SET IN THE BLOCK DATA. 4240
C 4250
C THE COMMON BLOCK /GEAR0/ CONTAINS THE VARIABLES DTUSED, NQUSED, 4260
C NSTEP, NFE, AND NJE AND CAN BE ACCESSED EXTERNALLY BY THE USER IF 4270
C DESIRED. RESPECTIVELY, IT CONTAINS THE STEP SIZE LAST USED (SUCCESS- 4280
C FULLY), THE ORDER LAST USED (SUCCESSFULLY), THE NUMBER OF STEPS TAKEN 4290
C SO FAR, THE NUMBER OF RESIDUAL EVALUATIONS (RES CALLS) SO FAR, 4300
C AND THE NUMBER OF MATRIX EVALUATIONS (PSETIB CALLS) SO FAR. 4310
C DIFFUN CALLS ARE COUNTED IN WITH RESIDUAL EVALUATIONS. 4320
C 4330
C THE OUTPUT PARAMETERS ARE.. 4340
C DT = THE STEP SIZE USED LAST, WHETHER SUCCESSFULLY OR NOT. 4350
C TOUT = THE OUTPUT VALUE OF T. IF INTEGRATION WAS SUCCESSFUL, 4360

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C AND THE INPUT VALUE OF INDEX WAS NOT 3, TOUT IS 4370
C UNCHANGED FROM ITS INPUT VALUE. OTHERWISE, TOUT 4380
C IS THE CURRENT VALUE OF T TO WHICH THE INTEGRATION 4390
C HAS BEEN COMPLETED. 4400
C INDEX = INTEGER USED ON OUTPUT TO INDICATE RESULTS, 4410
C WITH THE FOLLOWING VALUES AND MEANINGS.. 4420
C 0 INTEGRATION WAS COMPLETED TO TOUT OR BEYOND. 4430
C -1 THE INTEGRATION WAS HALTED AFTER FAILING TO PASS THE 4440
C ERROR TEST EVEN AFTER REDUCING DT BY A FACTOR OF 4450
C 1.E10 FROM ITS INITIAL VALUE. 4460
C -2 AFTER SOME INITIAL SUCCESS, THE INTEGRATION WAS 4470
C HALTED EITHER BY REPEATED ERROR TEST FAILURES OR BY 4480
C A TEST ON EPS. TOO MUCH ACCURACY HAS BEEN REQUESTED. 4490
C -3 THE INTEGRATION WAS HALTED AFTER FAILING TO ACHIEVE 4500
C CORRECTOR CONVERGENCE EVEN AFTER REDUCING DT BY A 4510
C FACTOR OF 1.E10 FROM ITS INITIAL VALUE. 4520
C -4 SINGULAR MATRIX ENCOUNTERED. PROBABLY DUE TO STORAGE 4530
C OVERWRITES. 4540
C -5 INDEX WAS 4 ON INPUT, BUT THE DESIRED CHANGES OF 4550
C PARAMETERS WERE NOT IMPLEMENTED BECAUSE TOUT 4560
C WAS NOT BEYOND T. INTERPOLATION TO T = TOUT WAS 4570
C PERFORMED AS ON A NORMAL RETURN. TO TRY AGAIN, 4580
C SIMPLY CALL AGAIN WITH INDEX = 4 AND A NEW TOUT. 4590
C -6 ILLEGAL INDEX VALUE. 4600
C -7 ILLEGAL EPS VALUE. 4610
C -8 AN ATTEMPT TO INTEGRATE IN THE WRONG DIRECTION. THE 4620
C SIGN OF DT IS WRONG RELATIVE TO T0 AND TOUT. 4630
C -9 DT .EQ. 0.0. 4640
C -10 ILLEGAL NINT VALUE. 4650
C -11 ILLEGAL KORD VALUE. 4660
C -12 ILLEGAL NCC VALUE. 4670
C -13 ILLEGAL NPDE VALUE. 4680
C -14 ILLEGAL MF VALUE. 4690
C -15 ILLEGAL BREAKPOINTS - NOT STRICTLY INCREASING. 4700
C -16 INSUFFICIENT STORAGE FOR WORK OR IWORK. 4710

SUBROUTINE VALUES (X, USOL, SCTCH, NDIM1, NDIM2, NPTS, NDERV, WORK)

----- 9250
C----- 9260
C SUBROUTINE VALUES COMPUTES THE SOLUTION U AND THE FIRST NDERV 9270
C DERIVATIVES OF U AT THE NPTS POINTS X AND AT TIME TOUT AND RETURNS 9280
C THEM IN THE ARRAY USOL. THIS ROUTINE MUST BE USED TO OBTAIN 9290
C SOLUTION VALUES SINCE PDECOL DOES NOT RETURN ANY SOLUTION VALUES 9300
C TO THE USER. SEE PDECOL. 9310
C 9320
C THE CALLING PARAMETERS ARE... 9330
C X = AN ARBITRARY VECTOR OF SPATIAL POINTS OF LENGTH NPTS AT 9340
C WHICH THE SOLUTION AND THE FIRST NDERV DERIVATIVE VALUES 9350
C ARE TO BE CALCULATED. IF X .LT. XLEFT (X .GT. XRIGHT) 9360
C THEN THE PIECEWISE POLYNOMIAL OVER THE LEFTMOST (RIGHT- 9370
C MOST) INTERVAL IS EVALUATED TO CALCULATE THE SOLUTION 9380
C VALUES AT THIS UNUSUAL VALUE OF X. SEE PDECOL. 9390
C 9400
C USOL = AN ARRAY WHICH CONTAINS THE SOLUTION AND THE FIRST 9410
C NDERV DERIVATIVES OF THE SOLUTION AT ALL THE POINTS IN 9420
C THE INPUT VECTOR X. IN PARTICULAR, USOL(J,I,K) CONTAINS 9430
C THE VALUE OF THE (K-1)-ST DERIVATIVE OF THE J-TH PDE 9440
C COMPONENT AT THE I-TH POINT OF THE X VECTOR FOR 9450
C J = 1 TO NPDE, I = 1 TO NPTS, AND K = 1 TO NDERV+1. 9460
C 9470
C SCTCH = A USER SUPPLIED WORKING STORAGE ARRAY OF LENGTH AT LEAST 9480
C KORD*(NDERV+1). SEE BELOW AND PDECOL FOR DEFINITIONS OF 9490
C THESE PARAMETERS. 9500
C 9510
C NDIM1 = THE FIRST DIMENSION OF THE OUTPUT ARRAY USOL IN THE CALLING 9520
C PROGRAM. NDIM1 MUST BE .GE. NPDE. 9530
C 9540
C NDIM2 = THE SECOND DIMENSION OF THE OUTPUT ARRAY USOL IN THE 9550
C CALLING PROGRAM. NDIM2 MUST BE .GE. NPTS. 9560
C 9570
C NPTS = THE NUMBER OF POINTS IN THE X VECTOR. 9580
C 9590
C NDERV = THE NUMBER OF DERIVATIVE VALUES OF THE SOLUTION THAT ARE 9600
C TO BE CALCULATED. NDERV SHOULD BE LESS THAN KORD SINCE 9610

C	THE KORD-TH DERIVATIVE OF A POLYNOMIAL OF DEGREE KORD-1	9620
C	IS EQUAL TO ZERO. SEE PDECOL.	9630
C		9640
C	WORK = THE USERS WORKING STORAGE ARRAY WHICH IS USED IN THIS CASE	9650
C	TO PROVIDE THE CURRENT BASIS FUNCTION COEFFICIENTS AND THE	9660
C	PIECEWISE POLYNOMIAL BREAKPOINT SEQUENCE.	9670
C		9680
C	PACKAGE ROUTINES CALLED.. BSPLVD,INTERV	9690
C	USER ROUTINES CALLED.. NONE	9700
C	CALLED BY.. USERS MAIN PROGRAM	9710
C	FORTRAN FUNCTIONS USED.. NONE	9720
C		9730
C	-----	9740
	BLOCK DATA	9950
C	-----	9960
C	IN THE FOLLOWING DATA STATEMENT, SET..	9970
C	LOUT = THE LOGICAL UNIT NUMBER FOR THE OUTPUT OF MESSAGES DURING	9980
C	THE INTEGRATION.	9990
C	NOGAUS = SET .EQ. 1 IF THE GAUSS-LEGENDRE COLLOCATION POINTS ARE	10000
C	NOT DESIRED WHEN NCC = 2 (SEE PDECOL AND COLPNT).	10010
C	MAXDER = SET .EQ. 5. ITS VALUE REPRESENTS THE MAXIMUM ORDER OF	10020
C	THE TIME INTEGRATION ALLOWED. ITS VALUE AFFECTS THE STOR-	10030
C	AGE REQUIRED IN WORK AND MAY BE CHANGED IF DESIRED	10040
C	(SEE COSET FOR RESTRICTIONS).	10050
C	UROUND = THE UNIT ROUNDOFF OF THE MACHINE, I.E. THE SMALLEST	10060
C	POSITIVE U SUCH THAT 1. + U .NE. 1. ON THE MACHINE.	10070
C	-----	10080
	COMMON /GEAR1/ DUM(5),UROUND, IDUM(4)	10090
	COMMON /OPTION/ NOGAUS,MAXDER	10100
	COMMON /IUNIT/ LOUT	10110
	DATA LOUT,NOGAUS,MAXDER,UROUND/3,0,5,7.1E-15/	10120
	END	10130

ALGORITHM 541

Efficient Fortran Subprograms for the Solution of Separable Elliptic Partial Differential Equations [D3]

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Language: Fortran

DESCRIPTION

1. Introduction

Computer models of geophysical processes often require the numerical solution of elliptic partial differential equations. This is particularly true for models that make use of stream functions, velocity potentials, or vorticity equations and for models that compute the pressure of an incompressible fluid. The numerical solution of elliptic equations can be a formidable programming task. Moreover, the equations are often time dependent, requiring repeated solutions and, hence, considerable computing resources.

Recent advances in computing methods [1, 2] made it possible to solve a very large class of elliptic equations (the separable ones) rapidly and with minimal storage. And as a result of work on singular problems [6, 8], this class is free of special cases for which solutions cannot be obtained numerically. This paper describes a package of computer programs that make use of current methods for solving elliptic partial differential equations. The package is fully documented in [7].

We first became involved in implementing the Buneman algorithm [2] and its extensions via the capacitance matrix approach [1] for solving Poisson's equation in polar coordinates on a disk. This led directly to two important problems: the need to treat the coordinate singularity which occurs at the center of the disk [8] and the restriction that the number of unknowns in one coordinate must be a power of 2 [10]. It is desirable to relax that restriction for certain studies, since geophysicists often work on grids that result in a number of points which are not

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a power of 2; for example, the commonly used 5° grid results in 72 grid points in longitude.

This work demonstrated that a Poisson equation in any of the usual coordinate systems was readily solvable, except on the interior of the sphere. The equation for that problem is a special case of the general separable equation (see Section 2, subroutine S.7) and cannot be solved by the Buneman algorithm. This difficulty led to the development of a cyclic reduction algorithm for separable elliptic equations [5]. We then wrote two subroutines, POIS and BLKTRI, to implement the cyclic reduction and generalized cyclic reduction, respectively. But because users were significantly inconvenienced by having to set up the finite-difference approximations before using either of these subroutines, the routines were generally *not* well accepted. Hence the National Center for Atmospheric Research (NCAR) supported the development of the driver programs S.1 through S.5 (listed in Section 2), which automated the process of generating the finite-difference equations for Poisson's equation in various coordinate systems. After these seven subroutines had been developed, we wrote a technical note [7] describing their use in detail. The combination of the ease of use of the subroutines and the complete documentation substantially increased the usage of the subroutines.

This experience made it quite apparent that even if a numerical method is excellent, the programs will not be widely used unless the software is easy to use and the documentation is quite clear. In addition to supporting user convenience, our philosophy has been to avoid terminology that might limit the class of users. For example, terms such as Dirichlet and Neumann boundary conditions, which tend to limit readership, have been replaced with either the solution or the derivative of the solution being specified.

The NCAR technical note [7] contains seven chapters, each describing one of the seven Fortran subroutines, and an Appendix dealing with least-squares solution of singular linear systems of equations. The first five subroutines solve a Helmholtz equation in Cartesian, polar, cylindrical, interior spherical, and surface spherical coordinates, respectively. The equations are solved on regular domains, namely, on a rectangle in the particular coordinate system that is chosen. For example, in polar coordinates the user can select the domain $a \leq r \leq b$ and $c \leq \theta \leq d$, where a can be zero. The user can also choose any of the standard boundary conditions. The solution or its derivative can be specified on the boundary, or periodic boundary conditions can be specified. Although these programs solve only two-dimensional problems, they can be adapted for use in three-dimensional problems by Fourier transforming in the third variable and using the parameter in the Helmholtz equation. This procedure is described in detail in the technical note [7]. Each chapter also contains a sample Fortran program which illustrates the use of the subroutine and provides a check during program installation.

In addition to the description of the Fortran subroutines, the first five chapters contain the finite-difference approximations used for the Helmholtz equation and the boundary conditions. Any special equations at singular points of the Helmholtz equation (e.g. at the origin, $r = 0$) are also presented. Approximations of second-order accuracy are used throughout. Each of these five chapters contains a section entitled "Singular Problems," which augments the Appendix in describing how solutions are obtained in the least-squares sense when a solution does not exist in the usual sense.

Each of these five subroutines calls one of the last two subroutines (POIS and BLKTRI), which can be used to solve a more general class of equation. Chapter VI describes subroutine POIS, which can be used to solve finite-difference approximations to an equation of the form

$$a(x) \frac{\partial^2 u}{\partial x^2} + b(x) \frac{\partial u}{\partial x} + c(x)u + \frac{\partial^2 u}{\partial y^2} = g(x, y).$$

Chapter VII describes subroutine BLKTRI, which can be used to solve finite-difference approximations to the general separable equation

$$a(x) \frac{\partial^2 u}{\partial x^2} + b(x) \frac{\partial u}{\partial x} + c(x)u + d(y) \frac{\partial^2 u}{\partial y^2} + e(y) \frac{\partial u}{\partial y} + f(y)u = g(x, y).$$

Although BLKTRI can solve problems that POIS solves, POIS is faster.

Subroutines POIS and BLKTRI use cyclic reduction and generalized cyclic reduction, respectively [2, 5, 10], to solve the large sparse linear systems. Although some discussion of these methods is included, program descriptions are presented separately for those who may wish to use the package only as a computational aid.

In Section 2 of this paper we briefly describe the capabilities of each of the seven subroutines; in Section 3 we summarize the various sections in the documentation; and in Section 4 we illustrate the documentation by presenting a small portion of the NCAR technical note [7]. And finally, in Section 5 we describe the modifications that have been made since the initial publication of the package and several developments that we are now planning or implementing.

2. Software Package

The package consists of seven subroutines. The first five, which are referred to as drivers, generate finite-difference approximations to a two-dimensional Helmholtz equation in a particular coordinate system. They are the following.

S.1 Subroutine PWSCRT solves the Helmholtz equation in Cartesian coordinates

$$\frac{\partial^2 u}{\partial x^2} + \frac{\partial^2 u}{\partial y^2} + \lambda u = f(x, y)$$

on the rectangle $a \leq x \leq b$, $c \leq y \leq d$.

S.2 Subroutine PWSPLR solves the Helmholtz equation in polar coordinates

$$\frac{1}{r} \frac{\partial}{\partial r} \left(r \frac{\partial u}{\partial r} \right) + \frac{1}{r^2} \frac{\partial^2 u}{\partial \theta^2} + \lambda u = f(r, \theta)$$

on the disk (or sector of the disk) $0 \leq a \leq r \leq b$, $0 \leq c \leq \theta \leq d \leq 2\pi$.

S.3 Subroutine PWSCYL solves the modified Helmholtz equation in cylindrical coordinates

$$\frac{1}{r} \frac{\partial}{\partial r} \left(r \frac{\partial u}{\partial r} \right) + \frac{\partial^2 u}{\partial z^2} + \frac{\lambda}{r^2} u = f(r, z)$$

on the slab of the cylinder $0 \leq a \leq r \leq b$, $c \leq z \leq d$.

S.4 Subroutine PWSSSP solves the Helmholtz equation in spherical surface coordinates

$$\frac{1}{\sin \theta} \frac{\partial}{\partial \theta} \left(\sin \theta \frac{\partial u}{\partial \theta} \right) + \frac{1}{\sin^2 \theta} \frac{\partial^2 u}{\partial \phi^2} + \lambda u = f(\theta, \phi)$$

on the sphere (or sector of the sphere) $0 \leq a \leq \theta \leq b \leq \pi$, $0 \leq c \leq \phi \leq d \leq 2\pi$.

S.5 Subroutine PWSCSP solves the modified Helmholtz equation in interior spherical coordinates

$$\frac{1}{r^2} \frac{\partial}{\partial r} \left(r^2 \frac{\partial u}{\partial r} \right) + \frac{1}{r^2 \sin \theta} \frac{\partial}{\partial \theta} \left(\sin \theta \frac{\partial u}{\partial \theta} \right) + \frac{\lambda}{r^2 \sin^2 \theta} u = f(\theta, r)$$

on all or a portion of a cross section of the sphere $0 \leq a \leq r \leq b$, $0 \leq c \leq \theta \leq d \leq \pi$.

On each of the boundaries, any of the following conditions can be specified.

B.1 The solution is specified (Dirichlet condition).

B.2 The derivative of the solution is specified (Neumann condition).

B.3 The solution is specified to be periodic.

B.4 The solution is unspecified (this condition applies when the domain of the equation includes a singularity of the coordinate system, e.g. when solving the Helmholtz equation in polar coordinates on the entire disk $0 \leq r \leq b$, $0 \leq \theta \leq 2\pi$).

Each of the five subroutines performs the following tasks.

T.1 defines the coefficients in the linear system of equations resulting from the finite-difference approximation.

T.2 incorporates the given boundary data into the right side of the equations.

T.3 perturbs the right side of the equation when the linear system is singular in order to guarantee that a solution exists.

Once the linear system has been defined, each of the routines calls one of the two following subroutines for solving linear systems.

S.6 Subroutine POIS solves systems resulting from finite-difference approximations to equations of the form

$$a(x) \frac{\partial^2 u}{\partial x^2} + b(x) \frac{\partial u}{\partial x} + c(x)u + \frac{\partial^2 u}{\partial y^2} = f(x, y).$$

S.7 Subroutine BLKTRI solves systems resulting from finite-difference approximations to general separable elliptic equations of the form

$$a(x) \frac{\partial^2 u}{\partial x^2} + b(x) \frac{\partial u}{\partial x} + c(x)u + d(y) \frac{\partial^2 u}{\partial y^2} + e(y) \frac{\partial u}{\partial y} + f(y)u = g(x, y).$$

3. Documentation of the Package

The documentation is completely contained in the NCAR technical note [7]. Each of the seven chapters is devoted to one subroutine and contains the following sections.

D.1 Calling Sequence and Parameter List Description. This section gives detailed descriptions of the parameters on input and output and information likely to be of value to a user deciding whether or not he can use the routine (e.g. memory space required, timing, and accuracy). This section will be of primary interest to those who wish to use the program strictly as a computational aid. This information is also contained on COMMENT cards at the beginning of each subroutine.

D.2 Difference Approximations. This section gives the detailed development of the difference approximations used for the equations, for the boundary conditions, and for the equations at coordinate singularities.

D.3 Singular Problems. This section gives the exact method used to perturb the right side of the equation when the user has specified a problem resulting in a singular linear system of equations.

D.4 Three-Dimensional Problems. This section demonstrates how to combine the use of the Fourier method with the routine in order to solve a Poisson equation in three space variables. (It is for this reason that the Helmholtz term involving λ sometimes has a function multiplying it.)

D.5 Sample Problem. This section gives a complete description of how to define input parameters for a selected differential equation. A Fortran implementation of the description is also given and certain output variables are given. These programs are included in the package as an aid to the user when the package is installed at his institution.

The documentation clearly states the locations of all machine-dependent constants used, e.g. π , and all required library routines, e.g. COS(X). The machine precision is required by BLKTRI, but it is computed (approximately) internally.

4. Sample Documentation

In this section we reproduce selected parts of the documentation for the routine PWSSSP to illustrate the use of the routine and the format of the documentation. We shall first state a continuous problem to be solved and then give a Fortran program that defines all the necessary input parameters. We then present the

formal parameter list description. Most of the following material is taken verbatim from Chapter V of the NCAR technical note [7], although in a different order.

Assume that we wish to approximate the solution of the Poisson equation

$$\frac{1}{\sin \theta} \frac{\partial}{\partial \theta} \left(\sin \theta \frac{\partial u}{\partial \theta} \right) + \frac{1}{\sin^2 \theta} \frac{\partial^2 u}{\partial \phi^2} = 2 - 6 \sin^2 \theta \sin^2 \phi$$

on the northern hemisphere $0 \leq \theta \leq \pi/2$, $0 \leq \phi \leq 2\pi$, subject to equatorial symmetry, i.e. $(\partial u / \partial \theta)(\pi/2, \phi) = 0$, $0 \leq \phi \leq 2\pi$.

The exact solution to this problem is

$$u(\theta, \phi) = \sin^2 \theta \sin^2 \phi + c,$$

where c is any constant. Hence, this is a singular problem.

We choose a 5° grid for our approximation. The Fortran program that defines the necessary input parameters for subroutine PWSSSP is given in the first program of the listings which follow.

```

C
C PROGRAM TO ILLUSTRATE THE USE OF PWSSSP
C
C DIMENSION F(19,73) ,BDTF(73) ,SINT(19) ,SINP(73) ,
1 W(647)
C
C THE VALUE OF IDIMF IS THE FIRST DIMENSION OF F. W IS DIMENSIONED
C 11*(M+1)+6X(N+1)=647 SINCE M=18 AND N=72
C
C PI = 3.141592653589793
C INTL = 0
C TS = 0
C TF = PI/2.
C M = 18
C MBDCND = 6
C PS = 0
C PF = PI+PI
C N = 72
C NBDKND = 0
C ELMBDA = 0.
C IDIMF = 19
C
C GENERATE SINES FOR USE IN SUBSEQUENT COMPUTATIONS
C
C DTHETA = TF/FLOAT(M)
C MPI = M+1
C DO 101 I=1,MPI
C SINT(I) = SIN(FLOAT(I-1)*DTHETA)
101 CONTINUE
C DPHI = (PI+PI)/FLOAT(N)
C NPI = N+1
C DO 102 J=1,NPI
C SINP(J) = SIN(FLOAT(J-1)*DPHI)
102 CONTINUE
C
C COMPUTE RIGHT SIDE OF EQUATION AND STORE IN F
C
C DO 104 J=1,NPI
C DO 103 I=1,MPI
C F(I,J) = 2.-6.*(SINT(I)*SINP(J))**2
103 CONTINUE
104 CONTINUE
C
C STORE DERIVATIVE DATA AT THE EQUATOR
C
C DO 105 J=1,NPI
C BDTF(J) = 0.
105 CONTINUE
C
C CALL PWSSSP (INTL,TS,TF,M,MBDCND,BDTS,BDTF,PS,PF,N,NBDKND,BDPS,
1 BDPF,ELMBDA,F,IDIMF,PERTRB,IERROR,W)
C
C COMPUTE DISCRETIZATION ERROR. SINCE PROBLEM IS SINGULAR, THE
C SOLUTION MUST BE NORMALIZED.

```

```

C
  ERR = 0
  DO 107 J=1,NP1
    DO 106 I=1,MP1
      Z = ABS(F(I,J)-(SINT(I)*SINP(J))**2-F(1,1))
      IF (Z .GT. ERR) ERR = Z
106   CONTINUE
107   CONTINUE
C
  PRINT 1001 , IERROR,ERR,PERTRB
  STOP
C
C
C
1001 FORMAT (///9H IERROR= ,I3,10X,22H DISCRETIZATION ERR = ,E12.4,
1         14H PERTURBATION=,E12.4)
C
  END

```

After the CALL to PWSSSP, the discretization error—the maximum absolute difference between the exact solution and the finite-difference approximation—is computed and printed along with the output parameter PERTRB. The printed output is

IERROR = 0

DISCRETIZATION ERR = 3.3811E - 03 PERTURBATION = 6.3220E - 04

The parameters in the call to PWSSSP are defined in the listing which appears at the end of this paper.

Accuracy and Timing. The execution time is proportional to $MN \log_2 N$ and is given in Table I for the NCAR Control Data 7600 computer. To test the accuracy of the method, a uniform random number generator was used to create an array $V(I, J)$, where $0 \leq V(I, J) \leq 1$. An array $F(I, J)$ was then computed by differencing $V(I, J)$ in double precision, using the difference equations that correspond to the sample problem. With $F(I, J)$ as a right side, subroutine PWSSSP was used to compute a solution $U(I, J)$. The maximum absolute difference between $U(I, J)$ and $V(I, J)$ is given in Table I. These errors can be compared with the accuracy of the NCAR 7600 (10^{-14}) to give a measure of the loss of significant digits due to roundoff. This error should not be confused with the discretization error given after the sample Fortran program.

5. Concluding Remarks

This package is not static, but is currently undergoing modifications to improve the solution techniques and extend the range of problems that can be solved. The following additions and modifications have already been made.

(1) Currently available is a subroutine, SEPELI, which discretizes the general separable elliptic equation. Thus the process of differencing the equations, incorporating the boundary conditions, and solving the resulting equations is automated as in the driver programs described in Section 2. Also included in this program is the option of obtaining fourth-order accurate solutions via the method of deferred corrections [3]. As in the other drivers, a weighted least-squares solution is obtained if that solution does not exist in the usual sense. More

Table I

M	N	Execution time (ms)	Maximum absolute difference
8	32	14	2.3×10^{-13}
16	64	60	1.3×10^{-12}
32	128	248	7.7×10^{-12}

Note: If NBDCND \neq 0, the execution times can be halved.

information may be obtained from John Adams, who has developed this program at NCAR.

(2) Replacements for the subroutines POIS and BLKTRI have been developed and tested and are in their final stage of documentation. The earlier versions of POIS and BLKTRI impose restrictions on the block order of the linear system of equations. The new algorithms eliminate this restriction. They also reduce to the efficient cyclic reduction algorithms when $N = 2^p$. The new algorithm for POIS is given in [9].

(3) With the use of this new algorithm, a subroutine, POISTG, has been developed [4] to solve finite-difference approximations to Helmholtz equations defined on staggered grids. This routine has been written in the format of POIS and BLKTRI. Drivers such as those described in Section 2 are now being written to solve finite-difference approximations to modified Helmholtz equations defined on a staggered grid.

In the early planning stage is a subroutine that will solve the general nonseparable linear elliptic equation in which the coefficients depend on both independent variables. However, it is anticipated that this program will use iterative rather than direct methods.

We feel that this package of high-quality, reliable, efficient programs for the solution of finite-difference approximations to elliptic equations is a significant step in the development of a complete package designed to solve an extremely wide class of elliptic partial differential equations.

NOTE ADDED IN PROOF. In the period since this paper was submitted for publication, all of the subroutines listed in modifications (1)–(3) above have been implemented in the latest version of the package. In addition, a driver for solving a three-dimensional Helmholtz equation in Cartesian coordinates, its associated solver, and two fast Fourier transform packages have been implemented. The interested reader should contact one of the authors at NCAR.

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9. SWEET, R.A. A cyclic reduction algorithm for block tridiagonal systems of arbitrary dimension. *SIAM J. Numer. Anal.* 14 (1977), 706–720.
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ALGORITHM

[Summary information and a part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

NAME(*n*): indicates a Fortran module with *n* records
 NAME^T(*n*): indicates "NAME" is included for testing purposes
 Contents: PWSCRT(416), PWSPLR(512), PWSCYL(455), PWSSSP(313),
 PWSSS1(356), PWSCSP(340), PWSCS1(247), POIS(153),
 POISGN(411), POINT(136), TRID(43), TPIDP(63),
 BLKTRI(200), BLKTRI(132), PROD(90), PRODP(107),
 CPROD(103), CPRODP(125), PPADD(119), PSGF(12),
 BSRH(17), PPSGF(9), PPSPF(9), COMPB(98), TQLRAT(162),
 NCHECK(29), APEXPS(9), STORE(5), MAIN1^T(101),
 MAIN2^T(101), MAIN3^T(114), MAIN4^T(95), MAIN5^T(140),
 MAIN6^T(110), MAIN7^T(117)

```

SUBROUTINE PWSSSP (INTL,TS,TF,M,MBDCND,BDTS,BDTF,PS,PF,N,NBDCND, POI01387
1          BDPS,BDPF,ELMBDA,F, IDIMF,PERTRB,IERROR,W) POI01388
C          POI01389
C SUBROUTINE PWSSSP SOLVES A FINITE DIFFERENCE APPROXIMATION TO THE POI01414
C HELMHOLTZ EQUATION IN SPHERICAL COORDINATES AND ON THE SURFACE OF POI01415
C THE UNIT SPHERE (RADIUS OF 1): POI01416
C          POI01417
C          (1/SIN(THETA))(D/DTHETA)(SIN(THETA)(D/DTHETA)U) POI01418
C          POI01419
C          + (1/SIN(THETA)**2)(D/DPHI)(D/DPHI)U POI01420
C          POI01421
C          + LAMBDA*U = F(THETA,PHI) POI01422
C          POI01423
C WHERE THETA IS COLATITUDE AND PHI IS LONGITUDE. POI01424
C          POI01425
C THE ARGUMENTS ARE DEFINED AS: POI01426
C          POI01427
C          * * * * * ON INPUT * * * * * POI01428
C          POI01429
C          POI01430
C INTL POI01431
C = 0 ON INITIAL ENTRY TO PWSSSP OR IF N,NBDCND,PS OR PF ARE POI01432
C CHANGED FROM A PREVIOUS CALL POI01433
C = 1 IF PS,PF,N AND NBDCND ARE UNCHANGED FROM A PREVIOUS CALL POI01434
C          POI01435
C NOTE: A CALL WITH INTL = 1 IS ABOUT 1 PERCENT FASTER THAN A POI01436
C CALL WITH INTL = 0 . POI01437
C          POI01438
C TS,TF POI01439
C THE RANGE OF THETA (COLATITUDE), I.E., TS .LE. THETA .LE. TF. POI01440
C TS MUST BE LESS THAN TF. TS AND TF ARE IN RADIANS. A TS OF POI01441
C ZERO CORRESPONDS TO THE NORTH POLE AND A TF OF PI CORRESPONDS TO POI01442
C THE SOUTH POLE. POI01443
C          POI01444
C M POI01445
C THE NUMBER OF PANELS INTO WHICH THE INTERVAL (TS,TF) IS POI01446
C SUBDIVIDED. HENCE, THERE WILL BE M+1 GRID POINTS IN THE POI01447
C THETA-DIRECTION GIVEN BY THETA(I) = (I-1)DTHETA+TS FOR POI01448
C I = 1,2,...,M+1, WHERE DTHETA = (TF-TS)/M IS THE PANEL WIDTH. POI01449
C          POI01450
C MBDCND POI01451
C INDICATES THE TYPE OF BOUNDARY CONDITION AT THETA = TS AND POI01452
C THETA = TF. POI01453
C          POI01454
C = 1 IF THE SOLUTION IS SPECIFIED AT THETA = TS AND THETA = TF. POI01455
C = 2 IF THE SOLUTION IS SPECIFIED AT THETA = TS AND THE POI01456
C DERIVATIVE OF THE SOLUTION WITH RESPECT TO THETA IS POI01457
C SPECIFIED AT THETA = TF (SEE NOTE 2 BELOW). POI01458
C = 3 IF THE DERIVATIVE OF THE SOLUTION WITH RESPECT TO THETA IS POI01459
C SPECIFIED AT THETA = TS AND THETA = TF (SEE NOTES 1,2 POI01460
C BELOW). POI01461
C = 4 IF THE DERIVATIVE OF THE SOLUTION WITH RESPECT TO THETA IS POI01462
C SPECIFIED AT THETA = TS (SEE NOTE 1 BELOW) AND THE POI01463
C SOLUTION IS SPECIFIED AT THETA = TF. POI01464
C = 5 IF THE SOLUTION IS UNSPECIFIED AT THETA = TS = 0 AND THE POI01465
C SOLUTION IS SPECIFIED AT THETA = TF. POI01466
C = 6 IF THE SOLUTION IS UNSPECIFIED AT THETA = TS = 0 AND THE POI01467
C DERIVATIVE OF THE SOLUTION WITH RESPECT TO THETA IS POI01468
C SPECIFIED AT THETA = TF (SEE NOTE 2 BELOW). POI01469
C = 7 IF THE SOLUTION IS SPECIFIED AT THETA = TS AND THE POI01470
C SOLUTION IS UNSPECIFIED AT THETA = TF = PI. POI01471

```

C = 8 IF THE DERIVATIVE OF THE SOLUTION WITH RESPECT TO THETA IS SPECIFIED AT THETA = TS (SEE NOTE 1 BELOW) AND THE SOLUTION IS UNSPECIFIED AT THETA = TF = PI. POI01472
 C POI01473
 C POI01474
 C = 9 IF THE SOLUTION IS UNSPECIFIED AT THETA = TS = \emptyset AND THETA = TF = PI. POI01475
 C POI01476
 C POI01477
 C NOTES: 1. IF TS = \emptyset , DO NOT USE MBDCND = 3,4, OR 8, BUT INSTEAD USE MBDCND = 5,6, OR 9 . POI01478
 C POI01479
 C 2. IF TF = PI, DO NOT USE MBDCND = 2,3, OR 6, BUT INSTEAD USE MBDCND = 7,8, OR 9 . POI01480
 C POI01481
 C POI01482
 C BDTs POI01483
 C A ONE-DIMENSIONAL ARRAY OF LENGTH N+1 THAT SPECIFIES THE VALUES OF THE DERIVATIVE OF THE SOLUTION WITH RESPECT TO THETA AT THETA = TS. WHEN MBDCND = 3,4, OR 8, POI01484
 C POI01485
 C POI01486
 C $BDTS(J) = (D/DTHETA)U(TS,PHI(J)), J = 1,2,\dots,N+1$. POI01487
 C POI01488
 C WHEN MBDCND HAS ANY OTHER VALUE, BDTs IS A DUMMY VARIABLE. POI01489
 C POI01490
 C BDTF POI01491
 C A ONE-DIMENSIONAL ARRAY OF LENGTH N+1 THAT SPECIFIES THE VALUES OF THE DERIVATIVE OF THE SOLUTION WITH RESPECT TO THETA AT THETA = TF. WHEN MBDCND = 2,3, OR 6, POI01492
 C POI01493
 C POI01494
 C $BDTF(J) = (D/DTHETA)U(TF,PHI(J)), J = 1,2,\dots,N+1$. POI01495
 C POI01496
 C POI01497
 C WHEN MBDCND HAS ANY OTHER VALUE, BDTF IS A DUMMY VARIABLE. POI01498
 C POI01499
 C PS,PF POI01500
 C THE RANGE OF PHI (LONGITUDE), I.E., PS .LE. PHI .LE. PF. PS MUST BE LESS THAN PF. PS AND PF ARE IN RADIANS. IF PS = \emptyset AND PF = 2*PI, PERIODIC BOUNDARY CONDITIONS ARE USUALLY PRESCRIBED. POI01501
 C POI01502
 C POI01503
 C POI01504
 C POI01505
 C N POI01506
 C THE NUMBER OF PANELS INTO WHICH THE INTERVAL (PS,PF) IS SUBDIVIDED. HENCE, THERE WILL BE N+1 GRID POINTS IN THE PHI-DIRECTION GIVEN BY $PHI(J) = (J-1)DPHI+PS$ FOR $J = 1,2,\dots,N+1$, WHERE $DPHI = (PF-PS)/N$ IS THE PANEL WIDTH. N MUST BE OF THE FORM $(2**P)(3**Q)(5**R)$ WHERE P, Q, AND R ARE ANY NON-NEGATIVE INTEGERS. N MUST BE GREATER THAN 2 . POI01507
 C POI01508
 C POI01509
 C POI01510
 C POI01511
 C POI01512
 C POI01513
 C NBDCND POI01514
 C INDICATES THE TYPE OF BOUNDARY CONDITION AT PHI = PS AND PHI = PF. POI01515
 C POI01516
 C POI01517
 C = \emptyset IF THE SOLUTION IS PERIODIC IN PHI, I.E., $U(I,1) = U(I,N+1)$. POI01518
 C POI01519
 C = 1 IF THE SOLUTION IS SPECIFIED AT PHI = PS AND PHI = PF (SEE NOTE BELOW). POI01520
 C POI01521
 C = 2 IF THE SOLUTION IS SPECIFIED AT PHI = PS (SEE NOTE BELOW) AND THE DERIVATIVE OF THE SOLUTION WITH RESPECT TO PHI IS SPECIFIED AT PHI = PF. POI01522
 C POI01523
 C POI01524
 C = 3 IF THE DERIVATIVE OF THE SOLUTION WITH RESPECT TO PHI IS SPECIFIED AT PHI = PS AND PHI = PF. POI01525
 C POI01526
 C = 4 IF THE DERIVATIVE OF THE SOLUTION WITH RESPECT TO PHI IS SPECIFIED AT PS AND THE SOLUTION IS SPECIFIED AT PHI = PF (SEE NOTE BELOW). POI01527
 C POI01528
 C POI01529
 C POI01530
 C NOTE: NBDCND = 1,2, OR 4 CANNOT BE USED WITH MBDCND = 5,6,7,8, OR 9 (THE FORMER INDICATES THAT THE SOLUTION IS SPECIFIED AT A POLE, THE LATTER INDICATES THAT THE SOLUTION IS UNSPECIFIED). USE INSTEAD POI01531
 C POI01532
 C POI01533
 C POI01534
 C POI01535
 C POI01536
 C POI01537
 C BDPS POI01538
 C A ONE-DIMENSIONAL ARRAY OF LENGTH M+1 THAT SPECIFIES THE VALUES OF THE DERIVATIVE OF THE SOLUTION WITH RESPECT TO PHI AT PHI = PS. WHEN NBDCND = 3 OR 4, POI01539
 C POI01540
 C POI01541
 C $BDPS(I) = (D/DPHI)U(THETA(I),PS), I = 1,2,\dots,M+1$. POI01542
 C POI01543
 C POI01544
 C WHEN NBDCND HAS ANY OTHER VALUE, BDPS IS A DUMMY VARIABLE. POI01545
 C POI01546
 C BDPF POI01547
 C A ONE-DIMENSIONAL ARRAY OF LENGTH M+1 THAT SPECIFIES THE VALUES POI01548

C OF THE DERIVATIVE OF THE SOLUTION WITH RESPECT TO PHI AT POI01549
 C PHI = PF. WHEN NBDKND = 2 OR 3, POI01550
 C POI01551
 C $BDPF(I) = (D/DPHI)U(THETA(I),PF)$, $I = 1,2,\dots,M+1$. POI01552
 C POI01553
 C WHEN NBDKND HAS ANY OTHER VALUE, BDPF IS A DUMMY VARIABLE. POI01554
 C POI01555
 C ELMBDA POI01556
 C THE CONSTANT LAMBDA IN THE HELMHOLTZ EQUATION. IF POI01557
 C LAMBDA .GT. 0, A SOLUTION MAY NOT EXIST. HOWEVER, PWSSSP WILL POI01558
 C ATTEMPT TO FIND A SOLUTION. POI01559
 C POI01560
 C F POI01561
 C A TWO-DIMENSIONAL ARRAY THAT SPECIFIES THE VALUE OF THE RIGHT POI01562
 C SIDE OF THE HELMHOLTZ EQUATION AND BOUNDARY VALUES (IF ANY). POI01563
 C FOR $I = 2,3,\dots,M$ AND $J = 2,3,\dots,N$ POI01564
 C POI01565
 C $F(I,J) = F(THETA(I),PHI(J))$. POI01566
 C POI01567
 C ON THE BOUNDARIES F IS DEFINED BY POI01568
 C POI01569

MBDKND	F(I,J)	F(M+1,J)	
-----	-----	-----	
1	U(TS,PHI(J))	U(TF,PHI(J))	
2	U(TS,PHI(J))	F(TF,PHI(J))	
3	F(TS,PHI(J))	F(TF,PHI(J))	
4	F(TS,PHI(J))	U(TF,PHI(J))	
5	F(0,PS)	U(TF,PHI(J))	$J = 1,2,\dots,N+1$
6	F(0,PS)	F(TF,PHI(J))	
7	U(TS,PHI(J))	F(PI,PS)	
8	F(TS,PHI(J))	F(PI,PS)	
9	F(0,PS)	F(PI,PS)	

C POI01570
 C POI01571
 C POI01572
 C POI01573
 C POI01574
 C POI01575
 C POI01576
 C POI01577
 C POI01578
 C POI01579
 C POI01580
 C POI01581
 C POI01582

NBDKND	F(I,1)	F(I,N+1)	
-----	-----	-----	
0	F(THETA(I),PS)	F(THETA(I),PS)	
1	U(THETA(I),PS)	U(THETA(I),PF)	
2	U(THETA(I),PS)	F(THETA(I),PF)	$I = 1,2,\dots,M+1$
3	F(THETA(I),PS)	F(THETA(I),PF)	
4	F(THETA(I),PS)	U(THETA(I),PF)	

C POI01583
 C POI01584
 C POI01585
 C POI01586
 C POI01587
 C POI01588
 C POI01589
 C POI01590
 C POI01591
 C F MUST BE DIMENSIONED AT LEAST $(M+1)*(N+1)$. POI01592
 C POI01593
 C NOTE POI01594
 C POI01595
 C IF THE TABLE CALLS FOR BOTH THE SOLUTION U AND THE RIGHT SIDE F POI01596
 C AT A CORNER THEN THE SOLUTION MUST BE SPECIFIED. POI01597
 C POI01598
 C IDIMF POI01599
 C THE ROW (OR FIRST) DIMENSION OF THE ARRAY F AS IT APPEARS IN THE POI01600
 C PROGRAM CALLING PWSSSP. THIS PARAMETER IS USED TO SPECIFY THE POI01601
 C VARIABLE DIMENSION OF F. IDIMF MUST BE AT LEAST M+1 . POI01602
 C POI01603
 C W POI01604
 C A ONE-DIMENSIONAL ARRAY THAT MUST BE PROVIDED BY THE USER FOR POI01605
 C WORK SPACE. THE LENGTH OF W MUST BE AT LEAST $11(M+1)+6(N+1)$. POI01606
 C POI01607
 C POI01608
 C * * * * * ON OUTPUT * * * * * POI01609
 C POI01610
 C F POI01611
 C CONTAINS THE SOLUTION U(I,J) OF THE FINITE DIFFERENCE POI01612
 C APPROXIMATION FOR THE GRID POINT (THETA(I),PHI(J)), POI01613
 C $I = 1,2,\dots,M+1$, $J = 1,2,\dots,N+1$. POI01614
 C POI01615
 C PERTRB POI01616
 C IF ONE SPECIFIES A COMBINATION OF PERIODIC, DERIVATIVE OR POI01617
 C UNSPECIFIED BOUNDARY CONDITIONS FOR A POISSON EQUATION POI01618
 C (LAMBDA = 0), A SOLUTION MAY NOT EXIST. PERTRB IS A CONSTANT, POI01619
 C CALCULATED AND SUBTRACTED FROM F, WHICH ENSURES THAT A SOLUTION POI01620
 C EXISTS. PWSSSP THEN COMPUTES THIS SOLUTION, WHICH IS A LEAST POI01621
 C SQUARES SOLUTION TO THE ORIGINAL APPROXIMATION. THIS SOLUTION POI01622
 C IS NOT UNIQUE AND IS UNNORMALIZED. THE VALUE OF PERTRB SHOULD POI01623
 C BE SMALL COMPARED TO THE RIGHT SIDE F. OTHERWISE , A SOLUTION POI01624
 C IS OBTAINED TO AN ESSENTIALLY DIFFERENT PROBLEM. THIS COMPARISON POI01625

C	SHOULD ALWAYS BE MADE TO INSURE THAT A MEANINGFUL SOLUTION HAS	POI01626
C	BEEN OBTAINED	POI01627
C		POI01628
C	IERROR	POI01629
C	AN ERROR FLAG THAT INDICATES INVALID INPUT PARAMETERS. EXCEPT	POI01630
C	FOR NUMBERS 0 AND 8, A SOLUTION IS NOT ATTEMPTED.	POI01631
C		POI01632
C	= 0 NO ERROR.	POI01633
C	= 1 TS .LT. 0 OR TF .GT. PI	POI01634
C	= 2 TS .GE. TF.	POI01635
C	= 3 MBDCND .LT. 1 OR MBDCND .GT. 9 .	POI01636
C	= 4 PS .GE. PF.	POI01637
C	= 5 N IS NOT OF THE FORM (2**P)(3**Q)(5**R) OR N .LE. 2 .	POI01638
C	= 6 NBDCND .LT. 0 OR NBDCND .GT. 4 .	POI01639
C	= 7 AN NBDCND OF 1,2, OR 4 IS USED WITH AN	POI01640
C	MBDCND OF 5,6,7,8, OR 9 .	POI01641
C	= 8 ELMBDA .GT. 0 .	POI01642
C	= 9 IDIMF .LT. M+1 .	POI01643
C	= 10 TS=0 AND MBDCND=3,4 OR 8 OR TF=PI AND MBDCND=2,3 OR 6	POI01644
C		POI01645
C	SINCE THIS IS THE ONLY MEANS OF INDICATING A POSSIBLY INCORRECT	POI01646
C	CALL TO PWSSSP, THE USER SHOULD TEST IERROR AFTER A CALL.	POI01647
C		POI01648
C	W	POI01649
C	CONTAINS INTERMEDIATE VALUES THAT MUST NOT BE DESTROYED IF	POI01650
C	PWSSSP WILL BE CALLED AGAIN WITH INTL = 1 .	POI01651

CERTIFICATION OF ALGORITHM 541

Efficient Fortran Subprograms for the Solution of Separable Elliptic Partial Differential Equations [D3]

[P.N. Swarztrauber and R.A. Sweet, *ACM Trans. Math. Software* 5, 3 (September 1979), 352-364.]

Michael Steuerwalt [Recd 10 Feb. 1977 and 31 Jan. 1979]

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INTRODUCTION

The problem of computing solutions to modified Helmholtz equations (or, more generally, separable linear elliptic equations) with simple boundary conditions on a rectangle in any of several coordinate frames arises frequently in applications and as an intermediate step in the solution of nonlinear and evolution problems. The cost of developing, documenting, and testing a package for this task is not negligible. Therefore several federal laboratories that can profitably use the National Center for Atmospheric Research (NCAR) package [2] chose to collaborate in its certification.

THE PACKAGE

Physically, the NCAR package consists of about 4700 lines of Fortran code (almost 40 percent are comments), 800 lines of example drivers, and 140 pages of documentation. The package cost about \$300,000 to develop. In comparison, EISPACK has 11,500 lines of code (49 percent are comments), 10,000 lines of example drivers, 551 pages of documentation [1], and cost about \$900,000.

The heart of the NCAR package is the two routines POIS and BLKTRI, which solve the linear systems arising from standard second-order finite-difference approximations of separable elliptic boundary value problems on rectangles. The package also includes five drivers that build the discrete system from the least possible information: the differential equation and boundary conditions, the geometric region, and the number of (evenly spaced) mesh points in each direction. Most users will communicate with the core routines only through these drivers.

THE CERTIFICATION EFFORT

Each of the five laboratories that had agreed to collaborate in the certification effort assumed responsibility for a particular driver:

PWSCRT: Air Force Weapons Laboratory, Kirtland Air Force Base, Albuquerque, N. Mex.

PWSPLR: Lawrence Livermore Laboratory, Livermore, Calif.

PWSCYL: Sandia Laboratories, Livermore, Calif.

PWSCSP: Sandia Laboratories, Albuquerque, N. Mex.

PWSSSP: Los Alamos Scientific Laboratory, Los Alamos, N. Mex.

Each laboratory agreed to

- (a) Compile the entire package.
- (b) Verify the results of the seven NCAR example programs.
- (c) Verify the correct working of the input error detection code.
- (d) For its particular driver, run a test problem using
 - (1) several permissible regions,
 - (2) several mesh sizes in each direction,
 - (3) all possible boundary conditions,
 - (4) zero and nonzero values of λ .
- (e) Evaluate the documentation.

It is important to note what we did *not* try to do:

(a) Explicitly test the core routines POIS and BLKTRI. Access to these routines was through the drivers alone.

(b) Make efficiency tests. The methods of the core routines are among the best direct methods available, but neither the authors nor the certifiers claim that there are no methods more efficient. In particular, we expect that higher-order discretization schemes would be more efficient for a given (small enough) accuracy.

(c) Make severe tests of the package's robustness.

The schedule outlined above entails a considerable effort. Among the five laboratories there are perhaps 10 different Fortran compilers, so the simple compilation of the package is a good test of its portability. The testing implied by (d) is substantial. For example, the driver PWSSSP admits 9 different possible boundary conditions in the θ direction and 5 in the ϕ direction; not all combinations are compatible, and some are valid only for certain geometries. To complete part (d) for the PWSSSP routine, 8 different regions were used with all possible valid boundary conditions, and with 5 different mesh sizes in the θ direction and 4 in the ϕ direction, for a total of 2360 runs per value of λ . Testing of the other drivers was similar. See the Appendix for details of the test problems.

RESULTS OF THE CERTIFICATION EFFORT

After initial tests, the certification team suggested several changes in the package and the documentation. All our recommendations were adopted. The most significant of these changes are the following.

(1) Version numbers for the package. The version we are certifying here is Version 2.

(2) Correction of initialization errors. The core routines POIS and BLKTRI perform some preliminary computations that need not be repeated if certain problem parameters remain unchanged. Our first tests had revealed program errors that could be avoided only by reinitializing every problem.

(3) Additional code to check for illegal combinations of problem parameters. The documentation clearly indicates that certain choices of boundary conditions are incompatible with particular geometries. For instance, in PWSSSP, three possible choices of the boundary condition at the final value TF of θ require that $TF = \pi$. A sample problem was run with these choices of boundary condition but with $TF < \pi$. The package did not check for such illegal combinations, but simply computed, producing wrong answers. In Version 2, the drivers will also set an error flag if, say, $|TF - \pi|$ is "too large." We feel that the package's robustness

is well enhanced by including this check, which has been done in a portable way and at little computing expense. On the other hand, we should remark that Swarztrauber and Sweet have indicated to us some uneasiness regarding the introduction of the imprecise and machine-dependent notion of "too large," and that no package, whatever its robustness or quality of documentation may be, can protect a user bent on self-immolation.

(4) An extensive list of errata and changes to the documentation. The documentation is neither so exhaustive as the EISPACK guide nor so rich in examples of the routines' use. This is no handicap: use of the NCAR routines is straightforward, whereas EISPACK often provides several options for doing a computation.

All the program tests were repeated for Version 2 to verify the corrections. The tests show that the methods are indeed of second-order accuracy. We deem the package to be valuable software of good quality. We especially commend its design, which permits users to communicate with it in familiar terms so that they do not have to grasp the mechanics of the discretization procedure. We believe the documentation will help that user who only wants answers to his problems to get those answers while remaining in blissful ignorance of details peripheral to his interests.

APPENDIX. DETAILS OF THE TEST PROBLEMS

Each driver may distinguish types of regions according to the geometry of the problem. Some combinations of boundary conditions and regions are not admissible. In this Appendix we tabulate the valid combinations for each driver; the entry "—" indicates no legal combination. We also list the intervals whose Cartesian products form the regions used in the certification tests, and give the true solution u of the various boundary value problems. The appropriate boundary values and the function f can be determined for a given problem from u , λ , and the region. For all problems the values of λ , M , and N used were

λ :	0.0	-1.0			
M :	9	18	36	72	144
N :	15	30	45	60	

PWSCRT

All 25 possible combinations of MBDCND and NBDCND are valid. The two regions used were

$$[A, B] = [-1/2, 1/2]$$

$$[C, D] = [-1/2, 1/2] \quad [0, 2].$$

The true solution was

$$u = \sin 2\pi(x + 1/8) \cos 2\pi(y + 1/8).$$

PWSPLR

This driver distinguishes two types of regions:

$$\begin{array}{l} 0 < A \\ 0 = A \end{array} \left\{ \begin{array}{l} P \\ Z \end{array} \right.$$

The compatibility table has 34 entries:

MBDCND	NBDCND				
	0	1	2	3	4
1	PZ	PZ	PZ	PZ	PZ
2	PZ	PZ	PZ	PZ	PZ
3	P	P	P	P	P
4	P	P	P	P	P
5	Z	—	—	Z	—
6	Z	—	—	Z	—

The four regions used were

$$[A, B] = [1/4, 1] \quad [0, 1]$$

$$[C, D] = [0, 2\pi/3] \quad [2\pi, 8\pi/3].$$

The true solution was

$$u = r^3 \cos(4\pi\theta/[D - C] + \pi/4).$$

PWSCYL

This driver distinguishes two types of regions:

$$0 < A \quad \begin{array}{|l} P \\ \hline Z \end{array}$$

$$0 = A$$

The compatibility table has 40 entries:

MBDCND	NBDCND				
	0	1	2	3	4
1	PZ	PZ	PZ	PZ	PZ
2	PZ	PZ	PZ	PZ	PZ
3	P	P	P	P	P
4	P	P	P	P	P
5	Z	Z	Z	Z	Z
6	Z	Z	Z	Z	Z

The two regions used were:

$$[A, B] = [1, 2] \quad [0, 1]$$

$$[C, D] = [0, 2\pi].$$

The true solution was

$$u = r^2 \cos(z + \pi/4).$$

PWSCSP

This driver recognizes eight types of regions:

$$0 < TS, TF < \pi \quad \begin{array}{|l} A \quad E \\ \hline B \quad F \\ C \quad G \\ D \quad H \end{array}$$

$$0 = TS, TF < \pi$$

$$0 < TS, TF = \pi$$

$$0 = TS, TF = \pi$$

The compatibility table has 72 entries:

MBDCND	NBDCND					
	1	2	3	4	5	6
1	ABCD	ABCD	ABCD	ABCD	—	—
2	AB	AB	AB	AB	—	—
3	A	A	A	A	E	E
4	AC	AC	AC	AC	—	—
5	BD	BD	BD	BD	—	—
6	B	B	B	B	F	F
7	CD	CD	CD	CD	—	—
8	C	C	C	C	G	G
9	D	D	D	D	H	H

The eight regions used were

$$[TS, TF] = [\pi/4, \pi/2] \quad [0, \pi/2] \quad [\pi/2, \pi] \quad [0, \pi]$$

$$[RS, RF] = [1, 2] \quad [0, 1].$$

The true solution was

$$u = r^4 \cos^4 \theta.$$

PWSSSP

This driver distinguishes four types of regions:

$0 < TS, TF < \pi$	A
$0 = TS, TF < \pi$	B
$0 < TS, TF = \pi$	C
$0 = TS, TF = \pi$	D

There are 59 entries in the compatibility table:

		NBDCND				
MBDCND	0	1	2	3	4	
1	ABCD	ABCD	ABCD	ABCD	ABCD	
2	AB	AB	AB	AB	AB	
3	A	A	A	A	A	
4	AC	AC	AC	AC	AC	
5	BD	—	—	BD	—	
6	B	—	—	B	—	
7	CD	—	—	CD	—	
8	C	—	—	C	—	
9	D	—	—	D	—	

The tests used eight regions:

$$\begin{aligned}
 [TS, TF] &= [\pi/4, \pi/2] & [0, \pi/2] & [\pi/2, \pi] & [0, \pi] \\
 [PS, PF] &= [\pi/2, 3\pi/2] & [0, 2\pi] & &
 \end{aligned}$$

The true solution was

$$u = \sin^2 \theta \cos 2\phi.$$

Table I. PWSCRT Test
 (Region = $[-0.5, 0.5] \times [-0.5, 0.5]$, $\lambda = 0.0$. Times are in milliseconds.)

M	N	MBC	NBC	Time		Error
9	15	0	0	8	0.224E-13	0.416E-01
18	30	0	0	32	0.120E-13	0.104E-01
36	45	0	0	96	0.971E-15	0.312E-02
72	60	0	0	271	0.205E-13	0.116E-02
9	15	0	1	6	0.0	0.299E-01
18	30	0	1	23	0.0	0.748E-02
36	45	0	1	71	0.0	0.224E-02
72	60	0	1	198	0.0	0.838E-03
9	15	1	0	6	0.0	0.299E-01
18	30	1	0	22	0.0	0.743E-02
36	45	1	0	65	0.0	0.226E-02
72	60	1	0	178	0.0	0.837E-03
9	15	1	1	4	0.0	0.346E-01
18	30	1	1	16	0.0	0.870E-02
36	45	1	1	49	0.0	0.263E-02
72	60	1	1	136	0.0	0.977E-03

Table II. PWSSSP Test
 (Region = $[0, \pi/2] \times [0, 2\pi]$, $\lambda = 0.0$. Times are in milliseconds.)

M	N	MBC	NBC	Time	PERTRB	Error
9	15	1	0	6	0.0	0.225E-01
18	30	1	0	22	0.0	0.572E-02
36	45	1	0	66	0.0	0.256E-02
72	60	1	0	180	0.0	0.145E-02
9	15	1	1	4	0.0	0.246E-01
18	30	1	1	16	0.0	0.631E-02
36	45	1	1	50	0.0	0.286E-02
72	60	2	1	137	0.0	0.162E-02
9	15	2	0	6	0.0	0.471E-01
18	30	2	0	23	0.0	0.120E-01
36	45	2	0	67	0.0	0.519E-02
72	60	2	0	181	0.0	0.287E-02
9	15	2	1	5	0.0	0.669E-01
18	30	2	1	17	0.0	0.170E-01
36	45	2	1	51	0.0	0.747E-02
72	60	2	1	138	0.0	0.414E-02

We close with some remarks about the accuracy and speed of the package. Tables I and II summarize some of the test results obtained on a 7600 at Los Alamos, using a local compiler similar to the FTN compiler (optimization level 2). The CPU times, given in milliseconds, reflect the influences of boundary condition and mesh size and are in very good agreement with times reported in [2]. Error is measured in the sup norm over the interior of the mesh of $(M + 1) \times (N + 1)$ points; the entries indicate that the finite-difference scheme of the package has second-order accuracy. To see this, note that if the scheme is of second order, then the error on the i th mesh is approximately

$$\text{error}_i = \alpha M_i^{-2} + \beta N_i^{-2},$$

where $M_i = 9 \cdot 2^i$, $N_i = 15 \cdot (i + 1)$, $0 \leq i \leq 3$. The Cartesian problem of Table I is symmetric about the line $x = y$, and so $\alpha = \beta$. We therefore expect the ratios $\text{error}_{i-1}/\text{error}_i$ to have the values of 4.00, 3.32, 2.69, respectively. For the spherical surface problem of Table II, α is nearly zero and the expected ratios are 4.00, 2.25, 1.78.

These observations—second-order accuracy and good agreement with the timings of [2]—hold true across the entire test set.

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ALGORITHM 542

Incomplete Gamma Functions [S14]

WALTER GAUTSCHI
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Key Words and Phrases: computation of incomplete gamma functions, Taylor's series, continued fractions
CR Categories: 5.12
Language: Fortran

DESCRIPTION

This algorithm implements the procedure developed in [1].

REFERENCES

- GAUTSCHI, W. A computational procedure for incomplete gamma functions. *ACM Trans. Math. Software* 5, 4 (Dec. 1979), 466-481.

ALGORITHM

```

      SUBROUTINE GAM(A, X, ACC, G, GSTAR, IFLG, IFLGST)      10
C     LET GAMMA(A) DENOTE THE GAMMA FUNCTION AND GAM(A,X) THE      20
C     (COMPLEMENTARY) INCOMPLETE GAMMA FUNCTION,                  30
C     GAM(A,X)=INTEGRAL FROM T=X TO T=INFINITY OF EXP(-T)*T**(A-1).  40
C     GAMSTAR(A,X) DENOTE TRICOMI:S FORM OF THE INCOMPLETE GAMMA  50
C     FUNCTION, WHICH FOR A.GT.0. IS DEFINED BY                   60
C     GAMSTAR(A,X)=(X**(-A)/GAMMA(A))*INTEGRAL FROM T=0 TO T=X OF  70
C     EXP(-T)*T**(A-1),                                           80
C     AND FOR A.LE.0. BY ANALYTIC CONTINUATION. FOR THE PURPOSE OF  90
C     THIS SUBROUTINE, THESE FUNCTIONS ARE NORMALIZED AS FOLLOWS& 100
C     GAM(A,X)/GAMMA(A), IF A.GT.0.,                               110
C     G(A,X)=                                                      120
C     EXP(X)*X**(-A)*GAM(A,X), IF A.LE.0.,                       130
C     GSTAR(A,X)=(X**A)*GAMSTAR(A,X).                             140
C     THE PROGRAM BELOW ATTEMPTS TO EVALUATE G(A,X) AND GSTAR(A,X), 150
C     BOTH TO AN ACCURACY OF ACC SIGNIFICANT DECIMAL DIGITS, FOR ARBI- 160
C     TRARY REAL VALUES OF A AND NONNEGATIVE VALUES OF X. THE SUB- 170
C     ROUTINE AUTOMATICALLY CHECKS FOR UNDERFLOW AND OVERFLOW CONDI- 180
C     TIONS AND RETURNS APPROPRIATE WARNINGS THROUGH THE OUTPUT PARA- 190
C     200
C     210
C     220
C     230
C     240
C     250
C     260
C     270
C     280

```

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C METERS IFLG, IFLGST. A RESULT THAT UNDERFLOWS IS RETURNED WITH	290
C THE VALUE 0., ONE THAT OVERFLOWS WITH THE VALUE OF THE LARGEST	300
C MACHINE-REPRESENTABLE NUMBER.	310
C	320
C NEAR LINES IN THE (A,X)-PLANE, A.LT.0., ALONG WHICH GSTAR	330
C VANISHES, THE ACCURACY SPECIFIED WILL BE ATTAINED ONLY IN TERMS	340
C OF ABSOLUTE ERROR, NOT RELATIVE ERROR. THERE ARE OTHER (RARE)	350
C INSTANCES IN WHICH THE ACCURACY ATTAINED IS SOMEWHAT LESS THAN	360
C THE ACCURACY SPECIFIED. THE DISCREPANCY, HOWEVER, SHOULD NEVER	370
C EXCEED ONE OR TWO (DECIMAL) ORDERS OF ACCURACY# NO INDICATION	380
C OF THIS IS GIVEN THROUGH ERROR FLAGS.	390
C	400
C	410
C PARAMETER LIST&	420
C	430
C A -- THE FIRST ARGUMENT OF G AND GSTAR. TYPE REAL.	440
C X -- THE SECOND ARGUMENT OF G AND GSTAR. TYPE REAL.	450
C ACC -- THE NUMBER OF CORRECT SIGNIFICANT DECIMAL DIGITS	460
C DESIRED IN THE RESULTS. TYPE REAL.	470
C G -- AN OUTPUT VARIABLE RETURNING THE VALUE OF G(A,X).	480
C TYPE REAL.	490
C GSTAR -- AN OUTPUT VARIABLE RETURNING THE VALUE OF	500
C GSTAR(A,X). TYPE REAL.	510
C IFLG -- AN ERROR FLAG INDICATING A NUMBER OF ERROR CONDI-	520
C TIONS IN G UPON EXECUTION. TYPE INTEGER.	530
C IFLGST -- AN ERROR FLAG INDICATING A NUMBER OF ERROR CONDI-	540
C TIONS IN GSTAR UPON EXECUTION. TYPE INTEGER.	550
C THE VALUES OF IFLG AND IFLGST HAVE THE FOLLOWING	560
C MEANINGS&	570
C 0 - NO ERROR CONDITION.	580
C 1 - ILLEGAL NEGATIVE ARGUMENT X. THE ROUTINE EXITS	590
C WITH THE VALUES ZERO FOR G AND GSTAR.	600
C 2 - INFINITELY LARGE RESULT AT X=0. THE ROUTINE	610
C RETURNS THE LARGEST MACHINE-REPRESENTABLE NUMBER.	620
C 3 - (ONLY FOR IFLGST) GSTAR IS INDETERMINATE AT	630
C A=0. AND X=0. THE ROUTINE RETURNS THE VALUE 1.,	640
C WHICH IS THE LIMIT VALUE AS X TENDS TO ZERO FOR	650
C FIXED A=0.	660
C 4 - THE RESULT UNDERFLOWS. IT IS SET EQUAL TO 0.	670
C 5 - THE RESULT OVERFLOWS. IT IS SET EQUAL TO THE	680
C LARGEST MACHINE-REPRESENTABLE NUMBER, WITH	690
C PROPER SIGN.	700
C 6 - CONVERGENCE FAILS WITHIN 600 ITERATIONS, EITHER	710
C IN TAYLOR:S SERIES OR IN LEGENDRE:S CONTINUED	720
C FRACTION. REASON UNKNOWN. THE COMPUTATION IS	730
C ABORTED AND THE ROUTINE RETURNS WITH ZERO	740
C VALUES FOR G AND GSTAR.	750
C	760
C ALL MACHINE-DEPENDENT PARAMETERS ARE COLLECTED IN THE FIRST	770
C DATA DECLARATION. THEY ARE AS FOLLOWS&	780
C	790
C PREC -- THE SINGLE PRECISION ACCURACY, TO BE SET APPROXI-	800
C MATELY EQUAL TO BETA*ALOG(2.)/ALOG(10.), WHERE BETA	810
C IS THE NUMBER OF BINARY DIGITS AVAILABLE IN THE MAN-	820
C TISSA OF THE SINGLE PRECISION FLOATING-POINT WORD.	830
C TYPE REAL.	840
C TOPEXP -- APPROXIMATELY THE LARGEST POSITIVE NUMBER T SUCH	850
C THAT 10.**T IS STILL REPRESENTABLE ON THE COMPUTER	860
C IN SINGLE PRECISION FLOATING-POINT ARITHMETIC.	870
C TYPE REAL.	880
C BOTEXP -- APPROXIMATELY THE SMALLEST NEGATIVE NUMBER T SUCH	890
C THAT 10.**T IS STILL REPRESENTABLE ON THE COMPUTER	900
C IN SINGLE PRECISION FLOATING-POINT ARITHMETIC.	910
C TYPE REAL.	920
C	930
C IN THE PROGRAM BELOW THESE PARAMETERS ARE SET TO CORRESPOND TO	940
C THE MACHINE CHARACTERISTICS OF THE CDC 6500 COMPUTER.	950
C	960
C THE SECOND DATA DECLARATION CONTAINS THE SINGLE PRECISION	970
C VALUE OF ALOG(10.). THE NEXT DATA DECLARATION CONTAINS THE SUCCES-	980
C SIVE COEFFICIENTS IN THE MACLAURIN EXPANSION OF (1/GAMMA(A+1))-1.	990
C THEY ARE GIVEN TO AS MANY DECIMAL PLACES AS IS NECESSARY TO ACHIEVE	1000
C MACHINE PRECISION (ON THE CDC 6500 COMPUTER) IN THE RANGE	1010
C ABS(A).LE..5. MORE ACCURATE VALUES OF THESE COEFFICIENTS (TO	1020
C 31 DECIMAL PLACES) CAN BE FOUND IN TABLE 5 OF J.W.WRENCH,JR.,	1030
C ::CONCERNING TWO SERIES FOR THE GAMMA FUNCTION::, MATH. COMPUT.	1040
C 22, 1968, 617-626.	

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C                                     1050
C THE SUBROUTINE CALLS ON A FUNCTION SUBROUTINE, NAMED ALGA, 1060
C WHICH IS TO PROVIDE SINGLE PRECISION VALUES OF THE LOGARITHM OF 1070
C THE GAMMA FUNCTION FOR ARGUMENTS LARGER THAN OR EQUAL TO .5. 1080
C A POSSIBLE VERSION OF SUCH A FUNCTION SUBROUTINE IS APPENDED 1090
C TO THE PRESENT SUBROUTINE. IT IS TAYLORED TO THE ACCURACY RE- 1100
C QUIREMENTS OF THE CDC 6500 COMPUTER, AND USES RATIONAL APPROXI- 1110
C MATIONS DUE TO CODY AND HILLSTROM (MATH. COMPUT. 21, 1967, 198- 1120
C 203). 1130
C 1140
C REFERENCE - W. GAUTSCHI, ::A COMPUTATIONAL PROCEDURE FOR 1150
C INCOMPLETE GAMMA FUNCTIONS::, ACM TRANS. MATH. SOFTWARE. 1160
C 1170
C DIMENSION C(18) 1180
C DATA PREC, TOPEXP, BOTEXP /1.445E1,3.22E2,-2.93E2/ 1190
C DATA AL10 /2.30258509299405E0/ 1200
C DATA C /.577215664901533E0,-.655878071520254E0, 1210
C * -4.2002635034095E-2,.16653861138229E0,-4.219773455554E-2, 1220
C * -9.62197152788E-3,7.21894324666E-3,-1.1651675919E-3, 1230
C * -2.152416741E-4,1.28050282E-4,-2.0134855E-5,-1.25049E-6, 1240
C * 1.13303E-6,-2.0563E-7,6.12E-9,5.00E-9,-1.2E-9,1.E-10/ 1250
C G = 0. 1260
C GSTAR = 0. 1270
C IF (X.LT.0.) GO TO 290 1280
C 1290
C INITIALIZATION 1300
C 1310
C IFLG = 0 1320
C IFLGST = 0 1330
C I = 0 1340
C IF (X.GT.0.) ALX = ALOG(X) 1350
C ALPHA = X + .25 1360
C IF (X.LT..25 .AND. X.GT.0.) ALPHA = ALOG(.5)/ALX 1370
C ALPREC = AL10*PREC 1380
C TOP = AL10*TOPEXP 1390
C BOT = AL10*BOTEXP 1400
C AINF = 10.**TOPEXP 1410
C EPS = .5*10.**(-ACC) 1420
C EPS1 = EPS/100. 1430
C SGA = 1. 1440
C IF (A.LT.0.) SGA = -SGA 1450
C AE = A 1460
C AA = ABS(A) 1470
C AP1 = A + 1. 1480
C AEPS1 = AP1 1490
C MA = .5 - A 1500
C FMA = FLOAT(MA) 1510
C AEPS = A + FMA 1520
C SGAE = 1. 1530
C IF (AEPS.LT.0.) SGAE = -SGAE 1540
C AAEPS = ABS(AEPS) 1550
C ALGP1 = 0. 1560
C 1570
C EVALUATION OF THE LOGARITHM OF THE ABSOLUTE VALUE OF 1580
C GAMMA(A+1.) AND DETERMINATION OF THE SIGN OF GAMMA(A+1.) 1590
C 1600
C SGGA = 1. 1610
C IF (MA.LE.0) GO TO 10 1620
C IF (AEPS.EQ.0.) GO TO 20 1630
C SGGA = SGAE 1640
C IF (MA.EQ.2*(MA/2)) SGGA = -SGGA 1650
C ALGP1 = ALGA(AEPS+1.) - ALOG(AEPS) 1660
C IF (MA.EQ.1) GO TO 20 1670
C ALGP1 = ALGP1 + ALGA(1.-AEPS) - ALGA(FMA-AEPS) 1680
C GO TO 20 1690
C 10 ALGP1 = ALGA(AP1) 1700
C 20 ALGEP1 = ALGP1 1710
C IF (X.GT.0.) GO TO 60 1720
C 1730
C EVALUATION OF GSTAR(A,0.) AND G(A,0.) 1740
C 1750
C IF (A) 30, 40, 50 1760
C 30 IFLGST = 2 1770
C GSTAR = AINF 1780
C G = 1./AA 1790
C RETURN 1800

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40	IFLGST = 3	1810
	GSTAR = 1.	1820
	IFLG = 2	1830
	G = AINF	1840
	RETURN	1850
50	G = 1.	1860
	RETURN	1870
60	IF (A.GT.ALPHA) GO TO 220	1880
	IF (X.GT.1.5) GO TO 240	1890
	IF (A.LT.-.5) GO TO 170	1900
C		1910
C	DIRECT EVALUATION OF G(A,X) AND GSTAR(A,X) FOR X.LE.1.5	1920
C	AND -.5.LE.A.LE.ALPHA(X)	1930
C		1940
	GSTAR = 1.	1950
	IF (A.GE..5) GO TO 110	1960
70	SUM = C(18)	1970
	DO 80 K=1,17	1980
	K1 = 18 - K	1990
	SUM = AE*SUM + C(K1)	2000
80	CONTINUE	2010
	GA = -SUM/(1.+AE*SUM)	2020
	Y = AE*ALX	2030
	IF (ABS(Y).GE.1.) GO TO 100	2040
	SUM = 1.	2050
	TERM = 1.	2060
	K = 1	2070
90	K = K + 1	2080
	IF (K.GT.600) GO TO 330	2090
	TERM = Y*TERM/FLOAT(K)	2100
	SUM = SUM + TERM	2110
	IF (ABS(TERM).GT.EPS1*SUM) GO TO 90	2120
	U = GA - SUM*ALX	2130
	GO TO 120	2140
100	U = GA - (EXP(Y)-1.)/AE	2150
	GO TO 120	2160
110	U = EXP(ALGA(A)) - (X**A)/A	2170
120	P = AE*X	2180
	Q = AE*P	2190
	R = AE + 3.	2200
	TERM = 1.	2210
	SUM = 1.	2220
	K = 1	2230
130	K = K + 1	2240
	IF (K.GT.600) GO TO 330	2250
	P = P + X	2260
	Q = Q + R	2270
	R = R + 2.	2280
	TERM = -P*TERM/Q	2290
	SUM = SUM + TERM	2300
	IF (ABS(TERM).GT.EPS1*SUM) GO TO 130	2310
	V = (X**AEP1)*SUM/AEP1	2320
	G = U + V	2330
	IF (I.EQ.1) GO TO 180	2340
	IF (A) 140, 150, 160	2350
140	T = EXP(X)*X**(-A)	2360
	G = T*G	2370
	GSTAR = 1. - A*G*EXP(-ALGP1)/T	2380
	RETURN	2390
150	G = EXP(X)*G	2400
	RETURN	2410
160	G = A*G*EXP(-ALGP1)	2420
	GSTAR = 1. - G	2430
	RETURN	2440
C		2450
C	RECURSIVE EVALUATION OF G(A,X) FOR X.LE.1.5 AND A.LT.-.5	2460
C		2470
170	I = 1	2480
	AE = AEPS	2490
	AEP1 = AEPS + 1.	2500
	IF (X.LT..25 .AND. AE.GT.ALPHA) GO TO 210	2510
	GO TO 70	2520
180	G = G*EXP(X)*X**(-AE)	2530
	DO 190 K=1,MA	2540
	G = (1.-X*G)/(FLOAT(K)-AE)	2550

190	CONTINUE	2560
	ALG = ALOG(G)	2570
C		2580
C	EVALUATION OF GSTAR(A,X) IN TERMS OF G(A,X)	2590
C		2600
200	GSSTAR = 1.	2610
	IF (MA.GE.0 .AND. AEPS.EQ.0.) RETURN	2620
	SGT = SGA*SGGA	2630
	T = ALOG(AA) - X + A*ALX + ALG - ALGP1	2640
	IF (T.LT.-ALPREC) RETURN	2650
	IF (T.GE.TOP) GO TO 320	2660
	GSSTAR = 1. - SGT*EXP(T)	2670
	RETURN	2680
210	I = 2	2690
	ALGP1 = ALGA(AEP1)	2700
C		2710
C	EVALUATION OF GSTAR(A,X) FOR A.GT.ALPHA(X) BY TAYLOR	2720
C	EXPANSION	2730
C		2740
220	G = 1.	2750
	TERM = 1.	2760
	SUM = 1.	2770
	K = 0	2780
230	K = K + 1	2790
	IF (K.GT.600) GO TO 340	2800
	TERM = X*TERM/(AE+FLOAT(K))	2810
	SUM = SUM + TERM	2820
	IF (ABS(TERM).GT.EPS*SUM) GO TO 230	2830
	ALGS = AE*ALX - X + ALOG(SUM) - ALGP1	2840
	IF (ALGS.LE.BOT) GO TO 310	2850
	GSSTAR = EXP(ALGS)	2860
	G = 1. - GSSTAR	2870
	IF (I.NE.2) RETURN	2880
	G = G*EXP(ALGP1)/AE	2890
	GO TO 180	2900
C		2910
C	EVALUATION OF G(A,X) FOR X.GT.1.5 AND A.LE.ALPHA(X) BY	2920
C	MEANS OF THE LEGENDRE CONTINUED FRACTION	2930
C		2940
240	GSSTAR = 1.	2950
	XPA = X + 1. - A	2960
	XMA = X - 1. - A	2970
	P = 0.	2980
	Q = XPA*XMA	2990
	R = 4.*XPA	3000
	S = -A + 1.	3010
	TERM = 1.	3020
	SUM = 1.	3030
	RHO = 0.	3040
	K = 1	3050
250	K = K + 1	3060
	IF (K.GT.600) GO TO 330	3070
	P = P + S	3080
	Q = Q + R	3090
	R = R + 8.	3100
	S = S + 2.	3110
	T = P*(1.+RHO)	3120
	RHO = T/(Q-T)	3130
	TERM = RHO*TERM	3140
	SUM = SUM + TERM	3150
	IF (ABS(TERM).GT.EPS*SUM) GO TO 250	3160
	IF (A) 260, 270, 280	3170
260	G = SUM/XPA	3180
	ALG = ALOG(G)	3190
	GO TO 200	3200
270	G = SUM/XPA	3210
	RETURN	3220
280	ALG = A*ALX - X + ALOG(A*SUM/XPA) - ALGP1	3230
	IF (ALG.LE.BOT) GO TO 300	3240
	G = EXP(ALG)	3250
	GSSTAR = 1. - G	3260
	RETURN	3270
290	IFLG = 1	3280
	IFLGST = 1	3290
	RETURN	3300
300	IFLG = 4	3310

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RETURN 3320
310 IFLGST = 4 3330
RETURN 3340
320 IFLGST = 5 3350
GSTAR = -SGT*AINF 3360
RETURN 3370
330 IFLG = 6 3380
RETURN 3390
340 IFLGST = 6 3400
RETURN 3410
END 3420

FUNCTION ALGA(X) 3430
DIMENSION CNUM(8), CDEN(8) 3440
DATA CNUM /4.120843185847770,85.68982062831317,243.175243524421, 3450
* -261.7218583856145,-922.2613728801522,-517.6383498023218, 3460
* -77.41064071332953,-2.208843997216182/, CDEN 3470
* /1.,45.64677187585908,377.8372484823942,951.323597679706, 3480
* 846.0755362020782,262.3083470269460,24.43519662506312, 3490
* .4097792921092615/ 3500
XI = AINT(X) 3510
IF (X-XI.GT..5) XI = XI + 1. 3520
M = IFIX(XI) - 1 3530
XE = X 3540
IF (M.EQ.-1) XE = X + 1. 3550
IF (M.GT.0) XE = X - FLOAT(M) 3560
SNUM = CNUM(1) 3570
SDEN = CDEN(1) 3580
DO 10 K=2,8 3590
SNUM = XE*SNUM + CNUM(K) 3600
SDEN = XE*SDEN + CDEN(K) 3610
10 CONTINUE 3620
ALGA = (XE-1.)*SNUM/SDEN 3630
IF (M.GT.-1) GO TO 20 3640
ALGA = ALGA - ALOG(X) 3650
RETURN 3660
20 IF (M.EQ.0) RETURN 3670
P = XE 3680
IF (M.EQ.1) GO TO 40 3690
MM1 = M - 1 3700
C 3710
C THE NEXT STATEMENT IS DESIGNED TO AVOID POSSIBLE OVERFLOW IN THE 3720
C COMPUTATION OF P. THE CONDITION IN THE IF-CLAUSE EXPRESSES THE 3730
C INEQUALITY  $1*3*5* \dots *(2*M+1)/(2**M).GE.Q$ , WHERE Q IS THE LARGEST 3740
C MACHINE-REPRESENTABLE NUMBER. 3750
C 3760
IF (M.GE.176) GO TO 50 3770
DO 30 K=1,MM1 3780
P = (XE+FLOAT(K))*P 3790
30 CONTINUE 3800
40 ALGA = ALGA + ALOG(P) 3810
RETURN 3820
50 ALGA = ALGA + ALOG(XE) 3830
DO 60 K=1,MM1 3840
ALGA = ALGA + ALOG(XE+FLOAT(K)) 3850
60 CONTINUE 3860
RETURN 3870
END 3880
C 3890
C DRIVER1 - ERROR FUNCTIONS 3900
C 3910
C ERF X FOR  $X=0(.05)1.5$  IN SINGLE AND DOUBLE PRECISION WITH 3920
C RELATIVE ERROR. CHECK AGAINST TABLE 7.1 IN NBS HANDBOOK. 3930
C 3940
DOUBLE PRECISION DPI, DC, DX, DXSQ, DG, DGSTAR, DERF, DXMSQ, DERFC 3950
PI = 4.*ATAN(1.) 3960
DPI = 4.D0*DATAN(1.D0) 3970
WRITE (6,99999) 3980
DO 10 I=1,31 3990
X = FLOAT(I-1)*.05 4000
DX = DBLE(FLOAT(I-1))*5.D-2 4010
XSQ = X*X 4020
DXSQ = DX*DX 4030
CALL GAM(.5, XSQ, 8., G, ERF, IFLG, IFLGST) 4040
CALL DGAM(.5D0, DXSQ, 16., DG, DERF, IFGD, IFGSTD) 4050

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        ERROR = SNGL(DABS(DBLE(ERF)-DERF))
        IF (DERF.NE.0.D0) ERROR = ABS(ERROR/SNGL(DERF))
        WRITE (6,99998) X, ERF, DERF, ERROR, IFLG, IFLGST, IFGD, IFGSTD
100 CONTINUE
C
C X*EXP(X*X)*ERFC X FOR X*(-2)=.005(.005).25 IN SINGLE AND
C DOUBLE PRECISION WITH RELATIVE ERROR. CHECK AGAINST TABLE 7.3
C IN NBS HANDBOOK.
C
        WRITE (6,99997)
        DO 30 I=1,50
            XMSQ = FLOAT(I)*.005
            DXMSQ = DBLE(FLOAT(I))*5.D-3
            XSQ = 1./XMSQ
            DXSQ = 1.D0/DXMSQ
            CALL GAM(.5, XSQ, 7., G, GSTAR, IFLG, IFLGST)
            CALL DGAM(.5D0, DXSQ, 16., DG, DGSTAR, IFGD, IFGSTD)
            ERF = 0.
            DERFC = 0.D0
            IF (XSQ.GT.740.) GO TO 20
            ERF = SQRT(XSQ)*EXP(XSQ)*G
            DERFC = DSQRT(DXSQ)*DEXP(DXSQ)*DG
200 ERROR = SNGL(DABS(DBLE(ERFC)-DERFC))
            IF (DERFC.NE.0.D0) ERROR = ABS(ERROR/SNGL(DERFC))
            WRITE (6,99996) XMSQ, ERF, DERFC, ERROR, IFLG, IFLGST, IFGD,
            * IFGSTD
300 CONTINUE
C
C ERF(SQRT(N*PI)) FOR N=1(1)10 IN SINGLE AND DOUBLE PRECISION
C WITH RELATIVE ERROR. CHECK AGAINST TABLE 7.3 IN NBS HANDBOOK.
C
        WRITE (6,99995)
        DO 40 N=1,10
            X = FLOAT(N)*PI
            DX = DBLE(FLOAT(N))*DPI
            CALL GAM(.5, X, 8., ERF, GSTAR, IFLG, IFLGST)
            CALL DGAM(.5D0, DX, 16., DERFC, DGSTAR, IFGD, IFGSTD)
            ERROR = ABS(SNGL((DBLE(ERFC)-DERFC)/DERFC))
            WRITE (6,99994) N, ERF, DERFC, ERROR, IFLG, IFLGST, IFGD,
            * IFGSTD
400 CONTINUE
        STOP
99999 FORMAT (/26X, 1HX, 8X, 5HERF X, 14X, 5HERF X, 12X, 5HERROR, 4X,
        * 4HIFLG, 1X, 6HIFLGST, 2X, 4HIFGD, 1X, 6HIFGSTD/)
99998 FORMAT (20X, E10.2, E15.7, D23.15, E10.2, 4I6)
99997 FORMAT (/23X, 7HX*(-2), 10X, 17HX*EXP(X*X)*ERFC X, 13X, 5HERROR,
        * 4X, 4HIFLG, 1X, 6HIFLGST, 2X, 4HIFGD, 1X, 6HIFGSTD/)
99996 FORMAT (20X, E10.2, E14.6, D23.15, E10.2, 4I6)
99995 FORMAT (/27X, 1HN, 13X, 16HERFC(SQRT(N*PI)), 14X, 5HERROR, 4X,
        * 4HIFLG, 1X, 6HIFLGST, 2X, 4HIFGD, 1X, 6HIFGSTD/)
99994 FORMAT (I28, E17.7, D23.15, E10.2, 4I6)
        END
C
C DRIVER3 - EXPONENTIAL INTEGRAL
C
C ESUBN(X) FOR N=0(1)20, X=VAR, IN SINGLE AND DOUBLE PRECISION
C
C ESUBN(X) FOR N=0(1)20, X=VAR, IN SINGLE AND DOUBLE PRECISION
C WITH RELATIVE ERROR. CHECK AGAINST TABLES I AND II IN PAGUROVA.
C
        DOUBLE PRECISION DX1, DX2, DX, DA, DG, DGSTAR, DESUBN, DP, DANU,
        * DESBNU, DLGA
        DIMENSION X1(10), X2(6), DX1(10), DX2(6)
        DATA X1 /0.,.01,.05,.2,.5,1.5,5.1,10.,14.7,19.8/
        DATA X2 /.01,.37,1.44,3.02,6.57,20./
        DATA DX1 /0.D0,1.D-2,5.D-2,.2D0,.5D0,1.5D0,5.1D0,1.D1,1.47D1,
        * 1.98D1/
        DATA DX2 /1.D-2,.37D0,1.44D0,3.02D0,6.57D0,2.D1/
        WRITE (6,99999)
        DO 50 I=1,10
            X = X1(I)
            DX = DX1(I)
            DO 40 J=1,21
                N = J - 1
                A = FLOAT(-N+1)

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      DA = DBLE(A) 4810
      CALL GAM(A, X, 8., G, GSTAR, IFLG, IFLGST) 4820
      CALL DGAM(DA, DX, 16., DG, DGSTAR, IFGD, IFGSTD) 4830
      IF (X.GT.0.) GO TO 10 4840
      ESUBN = 0. 4850
      DESUBN = 0.D0 4860
      IF (N.LE.1) GO TO 30 4870
      ESUBN = G 4880
      DESUBN = DG 4890
      GO TO 30 4900
10    IF (N.NE.0) GO TO 20 4910
      ESUBN = G/X 4920
      DESUBN = DG/DX 4930
      GO TO 30 4940
20    ESUBN = EXP(-X)*G 4950
      DESUBN = DEXP(-DX)*DG 4960
30    ERROR = SNGL(DABS(DBLE(ESUBN)-DESUBN)) 4970
      IF (DESUBN.NE.0.D0) ERROR = ABS(ERROR/SNGL(DESUBN)) 4980
      WRITE (6,99998) N, A, X, ESUBN, DESUBN, ERROR, IFLG, IFLGST, 4990
      * IFGD, IFGSTD 5000
40    CONTINUE 5010
      WRITE (6,99997) 5020
50    CONTINUE 5030
C 5040
C ESUBNU(X) FOR NU=0(.1)1., X=VAR, IN SINGLE AND DOUBLE 5050
C PRECISION WITH RELATIVE ERROR. CHECK AGAINST TABLE III IN 5060
C PAGUROVA. 5070
C 5080
      WRITE (6,99996) 5090
      DO 90 I=1,6 5100
      X = X2(I) 5110
      DX = DX2(I) 5120
      DO 80 J=1,11 5130
      ANU = FLOAT(J-1)*.1 5140
      A = -ANU + 1. 5150
      DANU = DBLE(FLOAT(J-1))*1.D0 5160
      DA = -DANU + 1.D0 5170
      CALL GAM(A, X, 7., G, GSTAR, IFLG, IFLGST) 5180
      CALL DGAM(DA, DX, 16., DG, DGSTAR, IFGD, IFGSTD) 5190
      IF (J.LT.11) GO TO 60 5200
      ESUBNU = EXP(-X)*G 5210
      DESBNU = DEXP(-DX)*DG 5220
      GO TO 70 5230
60    ESUBNU = X**(-A)*EXP(ALGA(A+1.))*G/A 5240
      DESBNU = DX**(-DA)*DEXP(DLGA(DA+1.D0))*DG/DA 5250
70    ERROR = SNGL(DABS(DBLE(ESUBNU)-DESBNU)) 5260
      IF (DESBNU.NE.0.D0) ERROR = ABS(ERROR/SNGL(DESBNU)) 5270
      WRITE (6,99995) ANU, A, X, ESUBNU, DESBNU, ERROR, IFLG, 5280
      * IFLGST, IFGD, IFGSTD 5290
80    CONTINUE 5300
      WRITE (6,99998) 5310
90    CONTINUE 5320
      STOP 5330
99999 FORMAT (/7X, 1HN, 5X, 1HA, 8X, 1HX, 8X, 8HESUBN(X), 10X, 6HESUBN(, 5340
      * 2HX), 10X, 5HERROR, 5X, 4HIFLG, 1X, 6HIFLGST, 2X, 4HIFGD, 1X, 5350
      * 6HIFGSTD/) 5360
99998 FORMAT (6X, I2, E10.1, E10.2, E15.7, D22.15, E10.2, 4I6) 5370
99997 FORMAT (/) 5380
99996 FORMAT (/5X, 2HNU, 7X, 1HA, 8X, 1HX, 8X, 9HESUBNU(X), 9X, 5390
      * 9HESUBNU(X), 9X, 5HERROR, 5X, 4HIFLG, 1X, 6HIFLGST, 2X, 4HIFGD, 5400
      * 1X, 6HIFGSTD/) 5410
99995 FORMAT (1X, 2E9.1, E10.2, E15.7, D22.15, E10.2, 4I6) 5420
      END 5430
C 5440
C DRIVER5 - CHISQUARE DISTRIBUTION 5450
C 5460
C P(CHISQUARE,NU) AND Q(CHISQUARE,NU) FOR SELECTED VALUES OF 5470
C CHISQUARE AND NU. CHECK AGAINST TABLE 26.7 IN NBS HANDBOOK. 5480
C 5490
      DIMENSION CCHSQ(9), NUMAX(9) 5500
      DATA CCHSQ /.1,1.,2.,4.,6.,8.,15.,20.,60./ 5510
      DATA NUMAX /6,12,16,22,27,4*30/ 5520
      WRITE (6,99999) 5530
      DO 20 I=1,9 5540
      CHSQ = CCHSQ(I) 5550
      NUI = NUMAX(I) 5560

```

```

DO 10 NU=1,NUI                                5570
  A = .5*FLOAT(NU)                             5580
  X = .5*CHSQ                                   5590
  CALL GAM(A, X, 5., Q, P, IFLG, IFLGST)        5600
  CALL GAM(A, X, 14., Q0, P0, IFLG0, IFLGS0)    5610
  ERRP = ABS(P-P0)/P0                          5620
  ERRQ = ABS(Q-Q0)/Q0                          5630
  WRITE (6,99998) NU, CHSQ, X, P, Q, ERRP, ERRQ, IFLG, IFLGST 5640
10 CONTINUE                                    5650
  WRITE (6,99997)                               5660
20 CONTINUE                                    5670
  STOP                                          5680
99999 FORMAT (//6X, 2HNU, 4X, 9HCHISQUARE, 7X, 1HX, 14X, 1HP, 14X, 1HQ, 5690
* 11X, 7HERROR P, 8X, 7HERROR Q, 5X, 4HIFLG, 2X, 6HIFLGST/) 5700
99998 FORMAT (5X, I3, 2E13.3, 4E15.4, 2I6)     5710
99997 FORMAT (/)                               5720
  END                                          5730

```

```

C                                              MAN7 10
C DRIVER7 - MOLECULAR INTEGRALS              MAN7 20
C ASUBN%ALPHA# FOR N#0%1F16 AND ALPHA#VAR TO 16 DECIMAL PLACES. MAN7 30
C CHECK AGAINST TABLES IN MILLER,GERHAUSEN AND MATSEN. MAN7 40
C                                              MAN7 50
  DOUBLE PRECISION ALPHA, X, P, A, G, GSTAR, DG, DGSTAR, ASUBN, MAN7 60
* DASUBN                                     MAN7 70
  DIMENSION ALPHA%10F                       MAN7 80
  DATA ALPHA /.125D0, .5D0, 1.625D0, 4.25D0, 7.375D0, 9.875D0, 12.625D0, MAN7 90
* 17.125D0, 21.25D0, 25.D0/                MAN7 100
  ACC # 16.                                  MAN7 110
  DACC # 25.                                 MAN7 120
  DO 40 I#1,10                               MAN7 130
    X # ALPHA%IF                             MAN7 140
    WRITE %6,99999F X                         MAN7 150
    WRITE %6,99998F                           MAN7 160
    DO 30 J#1,17                             MAN7 170
      N # J - 1                               MAN7 180
      P # 1.D0/X                             MAN7 190
      IF %N.EQ.0F GO TO 20                   MAN7 200
      DO 10 K#1,N                             MAN7 210
        P # DBLE%FLOAT%KFF*P/X              MAN7 220
10 CONTINUE                                  MAN7 230
20 A # DBLE%FLOAT%JFF                        MAN7 240
  CALL DGAM%A, X, ACC, G, GSTAR, IFLG, IFLGSTF MAN7 250
  CALL DGAM%A, X, DACC, DG, DGSTAR, IFLGD, IFLGSDF MAN7 260
  ASUBN # P*G                                MAN7 270
  DASUBN # P*DG                              MAN7 280
  ERR # ABS%SNGL%%ASUBN-DASUBN/DASUBNFF     MAN7 290
  WRITE %6,99997F N, ASUBN, DASUBN, ERR, IFLG, IFLGST, IFLGD, MAN7 300
* IFLGSD                                     MAN7 310
30 CONTINUE                                  MAN7 320
40 CONTINUE                                  MAN7 330
  STOP                                       MAN7 340
99999 FORMAT %/5X, 6HALPHA#, D11.4//F        MAN7 350
99998 FORMAT %12X, 1HN, 7X, 12HASUBN%ALPHA#F, 17X, 12HASUBN%ALPHA#F, 16X, MAN7 360
* 5HERROR, 6X, 4HIFLG, 1X, 6HIFLGST, 1X, 5HIFLGD, 1X, 6HIFLGD/F MAN7 370
99997 FORMAT %10X, I3, D24.15, D33.24, E15.4, 4I6F MAN7 380
  END                                       MAN7 390

```

```

SUBROUTINE DGAM(A, X, ACC, G, GSTAR, IFLG, IFLGST) 5740
  DOUBLE PRECISION A, X, G, GSTAR, C, ALL0, ALX, ALPHA, ALPREC, 5750
* TOP, BOT, AINF, EPS, EPS1, ES, SGA, AE, AA, AP1, AML, AEP1, 5760
* AEM1, FMA, AEPS, SGAE, AAEPS, ALGP1, SGG, ALGEP1, SGG, ALGS, 5770
* ALG, SUM, GA, Y, TERM, U, P, Q, R, V, T, H, SGT, ALX, RHO, XPA, 5780
* XMA, S, DLGA                               5790
  DIMENSION C(29)                             5800
  DATA PREC, TOPEXP, BOTEXP /28.8989, 322., -293./ 5810
  DATA ALL0 /2.3025850929940456840179914547D0/ 5820
  DATA C /.57721566490153286060651209008D0, -.65587807152025388107701 5830
* 951515D0, -4.200263503409523552900393488D-2, .166538611382291489501 5840
* 7007951D0, -4.21977345555443367482083013D-2, -9.6219715278769735621 5850
* 149217D-3, 7.2189432466630995423950103D-3, -1.165167591859065112113 5860
* 971D-3, -2.15241674114950972815730D-4, 1.2805028238811618615320D-4, 5870
* -2.013485478078823865569D-5, -1.2504934821426706573D-6, 5880

```

```

* 1.1330272319816958824D-6,-2.056338416977607103D-7,          5890
* 6.1160951044814158D-9,5.0020076444692229D-9,-1.181274570487020D-9  5900
* ,1.04342671169110D-10,7.782263439905D-12,-3.696805618642D-12,    5910
* 5.1003702875D-13,-2.058326054D-14,-5.34812254D-15,1.2267786D-15,  5920
* -1.181259D-16,1.187D-18,1.412D-18,-2.30D-19,1.7D-20/          5930
  G = 0.D0                                                         5940
  GSTAR = 0.D0                                                    5950
  IF (X.LT.0.D0) GO TO 290                                         5960
C                                                                     5970
C INITIALIZATION                                                  5980
C                                                                     5990
  IFLG = 0                                                         6000
  IFLGST = 0                                                       6010
  I = 0                                                             6020
  IF (X.GT.0.D0) ALX = DLOG(X)                                     6030
  ALPHA = X + .25D0                                                6040
  IF (X.LT..25D0 .AND. X.GT.0.D0) ALPHA = DLOG(.5D0)/ALX        6050
  ALPREC = AL10*DBLE(PREC)                                         6060
  TOP = AL10*DBLE(TOPEXP)                                         6070
  BOT = AL10*DBLE(BOTEXP)                                         6080
  AINF = 10.D0**TOPEXP                                           6090
  EPS = .5D0*10.D0**(-ACC)                                        6100
  EPS1 = EPS/1.D2                                                 6110
  SGA = 1.D0                                                       6120
  IF (A.LT.0.D0) SGA = -SGA                                       6130
  AE = A                                                            6140
  AA = DABS(A)                                                      6150
  AP1 = A + 1.D0                                                   6160
  AEP1 = AP1                                                       6170
  MA = SNGL(.5D0-A)                                               6180
  FMA = DBLE(FLOAT(MA))                                          6190
  AEPS = A + FMA                                                  6200
  SGAE = 1.D0                                                     6210
  IF (AEPS.LT.0.D0) SGAE = -SGAE                                  6220
  AAEPS = DABS(AEPS)                                              6230
  ALGP1 = 0.D0                                                    6240
C                                                                     6250
C EVALUATION OF THE LOGARITHM OF THE ABSOLUTE VALUE OF          6260
C GAMMA(A+1.) AND DETERMINATION OF THE SIGN OF GAMMA(A+1.)    6270
C                                                                     6280
  SGGA = 1.D0                                                     6290
  IF (MA.LE.0) GO TO 10                                           6300
  IF (AEPS.EQ.0.D0) GO TO 20                                       6310
  SGGA = SGAE                                                      6320
  IF (MA.EQ.2*(MA/2)) SGGA = -SGGA                               6330
  ALGP1 = DLGA(AEPS+1.D0) - DLOG(AEPS)                            6340
  IF (MA.EQ.1) GO TO 20                                           6350
  ALGP1 = ALGP1 + DLGA(1.D0-AEPS) - DLGA(FMA-AEPS)              6360
  GO TO 20                                                         6370
10 ALGP1 = DLGA(AP1)                                              6380
20 ALGEP1 = ALGP1                                                 6390
  IF (X.GT.0.D0) GO TO 60                                         6400
C                                                                     6410
C EVALUATION OF GSTAR(A,0.) AND G(A,0.)                          6420
C                                                                     6430
  IF (A) 30, 40, 50                                              6440
30 IFLGST = 2                                                     6450
  GSTAR = AINF                                                     6460
  G = 1.D0/AA                                                      6470
  RETURN                                                           6480
40 IFLGST = 3                                                     6490
  GSTAR = 1.D0                                                    6500
  IFLG = 2                                                         6510
  G = AINF                                                         6520
  RETURN                                                           6530
50 G = 1.D0                                                       6540
  RETURN                                                           6550
60 IF (A.GT.ALPHA) GO TO 220                                       6560
  IF (X.GT.1.5D0) GO TO 240                                       6570
  IF (A.LT.-.5D0) GO TO 170                                       6580
C                                                                     6590
C DIRECT EVALUATION OF G(A,X) AND GSTAR(A,X) FOR X.LE.1.5    6600
C AND -.5.LE.A.LE.ALPHA(X)                                       6610
C                                                                     6620
  GSTAR = 1.D0                                                    6630
  IF (A.GE..5D0) GO TO 110                                         6640

```

```

70 SUM = C(29)
DO 80 K=1,28
    K1 = 29 - K
    SUM = AE*SUM + C(K1)
80 CONTINUE
GA = -SUM/(1.D0+AE*SUM)
Y = AE*ALX
IF (DABS(Y).GE.1.D0) GO TO 100
SUM = 1.D0
TERM = 1.D0
K = 1
90 K = K + 1
IF (K.GT.600) GO TO 330
TERM = Y*TERM/DBLE(FLOAT(K))
SUM = SUM + TERM
IF (DABS(TERM).GT.EPS1*SUM) GO TO 90
U = GA - SUM*ALX
GO TO 120
100 U = GA - (DEXP(Y)-1.D0)/AE
GO TO 120
110 U = DEXP(DLGA(A)) - (X**A)/A
120 P = AE*X
Q = AEP1
R = AE + 3.D0
TERM = 1.D0
SUM = 1.D0
K = 1
130 K = K + 1
IF (K.GT.600) GO TO 330
P = P + X
Q = Q + R
R = R + 2.D0
TERM = -P*TERM/Q
SUM = SUM + TERM
IF (DABS(TERM).GT.EPS1*SUM) GO TO 130
V = (X**AEP1)*SUM/AEP1
G = U + V
IF (I.EQ.1) GO TO 180
IF (A) 140, 150, 160
140 T = DEXP(X)*X**(-A)
G = T*G
GSTAR = 1.D0 - A*G*DEXP(-ALGP1)/T
RETURN
150 G = DEXP(X)*G
RETURN
160 G = A*G*DEXP(-ALGP1)
GSTAR = 1.D0 - G
RETURN
C
C RECURSIVE EVALUATION OF G(A,X) FOR X.LE.1.5 AND A.LT.-.5
C
170 I = 1
AE = AEPS
AEP1 = AEPS + 1.D0
IF (X.LT..25D0 .AND. AE.GT.ALPHA) GO TO 210
GO TO 70
180 G = G*DEXP(X)*X**(-AE)
DO 190 K=1,MA
    G = (1.D0-X*G)/(DBLE(FLOAT(K))-AE)
190 CONTINUE
ALG = DLOG(G)
C
C EVALUATION OF GSTAR(A,X) IN TERMS OF G(A,X)
C
200 GSTAR = 1.D0
IF (MA.GE.0 .AND. AEPS.EQ.0.D0) RETURN
SGT = SGA*SGGA
T = DLOG(AA) - X + A*ALX + ALG - ALGP1
IF (T.LT.-ALPREC) RETURN
IF (T.GE.TOP) GO TO 320
GSTAR = 1.D0 - SGT*DEXP(T)
RETURN
210 I = 2
ALGP1 = DLGA(AEP1)
C
C EVALUATION OF GSTAR(A,X) FOR A.GT.ALPHA(X) BY TAYLOR

```

```

C EXPANSION 7410
C 7420
220 G = 1.D0 7430
    TERM = 1.D0 7440
    SUM = 1.D0 7450
    K = 0 7460
230 K = K + 1 7470
    IF (K.GT.600) GO TO 340 7480
    TERM = X*TERM/(AE+DBLE(FLOAT(K))) 7490
    SUM = SUM + TERM 7500
    IF (DABS(TERM).GT.EPS*SUM) GO TO 230 7510
    ALGS = AE*ALX - X + DLOG(SUM) - ALGEP1 7520
    IF (ALGS.LE.BOT) GO TO 310 7530
    GSTAR = DEXP(ALGS) 7540
    G = 1.D0 - GSTAR 7550
    IF (I.NE.2) RETURN 7560
    G = G*DEXP(ALGEP1)/AE 7570
    GO TO 180 7580
C 7590
C EVALUATION OF G(A,X) FOR X.GT.1.5 AND A.LE.ALPHA(X) BY 7600
C MEANS OF THE LEGENDRE CONTINUED FRACTION 7610
C 7620
240 GSTAR = 1.D0 7630
    XPA = X + 1.D0 - A 7640
    XMA = X - 1.D0 - A 7650
    P = 0.D0 7660
    Q = XPA*XMA 7670
    R = 4.D0*XPA 7680
    S = -A + 1.D0 7690
    TERM = 1.D0 7700
    SUM = 1.D0 7710
    RHO = 0.D0 7720
    K = 1 7730
250 K = K + 1 7740
    IF (K.GT.600) GO TO 330 7750
    P = P + S 7760
    Q = Q + R 7770
    R = R + 8.D0 7780
    S = S + 2.D0 7790
    RHO = T/(Q-T) 7810
    TERM = RHO*TERM 7820
    SUM = SUM + TERM 7830
    IF (DABS(TERM).GT.EPS*SUM) GO TO 250 7840
    IF (A) 260, 270, 280 7850
260 G = SUM/XPA 7860
    ALG = DLOG(G) 7870
    GO TO 200 7880
270 G = SUM/XPA 7890
    RETURN 7900
280 ALG = A*ALX - X + DLOG(A*SUM/XPA) - ALGPI 7910
    IF (ALG.LE.BOT) GO TO 300 7920
    G = DEXP(ALG) 7930
    GSTAR = 1.D0 - G 7940
    RETURN 7950
290 IFLG = 1 7960
    IFLGST = 1 7970
    RETURN 7980
300 IFLG = 4 7990
    RETURN 8000
310 IFLGST = 4 8010
    RETURN 8020
320 IFLGST = 5 8030
    GSTAR = -SGT*AINF 8040
    RETURN 8050
330 IFLG = 6 8060
    RETURN 8070
340 IFLGST = 6 8080
    RETURN 8090
    END 8100
    DOUBLE PRECISION FUNCTION DLGA(DX) 8110
    DOUBLE PRECISION DBNUM, DBDEN, DX, DC, DP, DY, DT, DS 8120
    DIMENSION DBNUM(8), DBDEN(8) 8130
    DATA DBNUM /-3.617D3,1.D0,-6.91D2,1.D0,-1.D0,1.D0,-1.D0,1.D0/, 8140
    * DBDEN /1.224D5,1.56D2,3.6036D5,1.188D3,1.68D3,1.26D3,3.6D2,1.2D1/ 8150
    DC = .5D0*DLOG(8.D0*DATAN(1.D0)) 8160
    DP = 1.D0 8170

```


COLLECTED ALGORITHMS (cont.)

542-P13- 0

DY = DX	8180
Y = SNGL(DY)	8190
C	8200
C THE CONDITIONAL CLAUSE IN THE NEXT STATEMENT EXPRESSES THE	8210
C INEQUALITY Y.GT.EXP(.121189*DPREC+.053905), WHERE DPREC IS THE	8220
C NUMBER OF DECIMAL DIGITS CARRIED IN DOUBLE PRECISION FLOATING-POINT	8230
C ARITHMETIC.	8240
C	8250
10 IF (Y.GT.35.027) GO TO 20	8260
DP = DY*DP	8270
DY = DY + 1.D0	8280
Y = SNGL(DY)	8290
GO TO 10	8300
20 DT = 1.D0/(DY*DY)	8310
DS = 4.3867D4/2.44188D5	8320
DO 30 I=1,8	8330
DS = DT*DS + DBNUM(I)/DBDEN(I)	8340
30 CONTINUE	8350
DLGA = (DY-.5D0)*DLOG(DY) - DY + DC + DS/DY - DLOG(DP)	8360
RETURN	8370
END	8380

ALGORITHM 543

FFT9, Fast Solution of Helmholtz-Type Partial Differential Equations [D3]

E. N. HOUSTIS
Purdue University
and
T. S. PAPATHEODOROU
Clarkson College of Technology

Key Words and Phrases: fast Fourier transform, fast Helmholtz solver, fast Poisson solver
CR Categories: 5.17
Language: Fortran

DESCRIPTION

The algorithm given here is a complement to [1] where the description, test results, and references are given.

REFERENCES

1. HOUSTIS, E.N., AND PAPATHEODOROU, T.S. High-order fast elliptic equation solver. *ACM Trans. Math. Software* 5, 4 (Dec. 1979), 431-441.

ALGORITHM

```

C      PROGRAM FFT9(INPUT,OUTPUT,TAPE5=INPUT,TAPE6=OUTPUT)          10
C                                                                 20
C      ---- PROGRAM DESCRIPTION ----                                30
C      PROGRAM FFT9 USES A 4-TH OR 6-TH ORDER 9-POINT DIFFERENCE    40
C      FORMULA AND FAST FOURIER TRANSFORM FOR THE NUMERICAL        50
C      SOLUTION OF THE ELLIPTIC EQUATION WITH CONSTANT COEFFICIENTS 60
C                                                                 70
C      (I) CUXX*DDXU + CUY Y*DDYU + CU*U = R                        80
C                                                                 90
C      ON A RECTANGULAR REGION 0 .LE. X .LE. SX,0 .LE. Y .LE. SY   100
C      AND SUBJECT TO DIRICHLET BOUNDARY CONDITIONS                110
C      (II) U = G                                                    120
C                                                                 130
C      NOTE- THE 6-TH ORDER ALGORITHM IS APPLIED ONLY             140
C      TO POISSON TYPE OPERATORS                                    150
C                                                                 160
C                                                                 170
C      ---- INPUT AND OUTPUT TO FFT9 ----                            180
C      --PROBLEM DEFINITION-- USER SUPPLIED FORTRAN FUNCTION FOR THE 190
C      EVALUATION OF THE RIGHT SIDES (R,G) OF THE DIFFERENTIAL     200

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```

C      AND BOUNDARY OPERATORS                                210
C      REAL          FUNCTION PDERGH(X,Y)                   220
C      REAL          FUNCTION PDERGH(X,Y)                   230
C      PDERGH = R                                           240
C      RETURN                                               250
C      END                                                  260
C      REAL          FUNCTION BCOND(I,X,Y,BVALUS)           270
C      TWO DIMENSIONS                                       280
C      VALUES OF BOUNDARY CONDITION ON SIDE I              290
C      AT (X,Y)                                             300
C      I=4                                                  310
C      -----                                             320
C      I              I                                     330
C      I              I                                     340
C      I=3 I REGION I I=1                                   350
C      I              I                                     360
C      I              I                                     370
C      I              I                                     380
C      -----                                             390
C      I=2                                                  400
C      REAL          BVALUS(4)                              410
C      GO TO(100,101,102,103) , I                          420
C      101 BVALUS(4) = G                                     430
C      BCOND = BVALUS(4)                                    440
C      RETURN                                               450
C      102 BVALUS(4) = G                                     460
C      BCOND = BVALUS(4)                                    470
C      RETURN                                               480
C      103 BVALUS(4) = G                                     490
C      BCOND = BVALUS(4)                                    500
C      RETURN                                               510
C      104 BVALUS(4) = G                                     520
C      BCOND = BVALUS(4)                                    530
C      RETURN                                               540
C      END                                                  550
C      USER SUPPLIED SUBROUTINE FOR THE DEFINITION OF     560
C      P.D.E CONSTANT COEFFICIENTS                         570
C      SUBROUTINE PDE(X,Y,CVALUS)                          580
C      REAL CVALUS(7)                                      590
C      CVALUS(1) = CUXX                                     600
C      CVALUS(3) = CUYX                                     610
C      CVALUS(6) = CU                                       620
C      RETURN                                               630
C      END                                                  640
C      USER SUPPLIED FORTRAN FUNCTION FOR THE TRUE SOLUTION 650
C      IF KNOWN                                            660
C      REAL          FUNCTION TRUE(X,Y)                     670
C      TRUE = ...                                          680
C      RETURN                                               690
C      END                                                  700
C      --REGION AND GRID SPECIFICATIONS--                  710
C      SX,SY          - LENGTHS OF SIDES OF RECTANGLE     720
C      NGRIDX,NGRIDY  - NUMBER OF HORIZONTAL AND          730
C      VERTICAL MESH LINES                                 740
C      IQX,IQY        - EXPONENTS OF 2                    750
C      NGRIDX=2**IQX+1,NGRIDY=2**IQY+1                    760
C      READ(5,100) SX,SY,NGRIDX,NGRIDY                     770
C      100 FORMAT(2F10.0,2I3)                               780
C      --OUTPUT CONTROL--USER SUPPLIED DATA              790
C      LEVEL          - OUTPUT LEVEL DESIRED               800
C      NRUNS          - NUMBER OF SUCCESIVE RUNS.          810
C      IN EACH RUN THE MESH SIZE IS                       820
C      CUT BY A FACTOR OF 2.                              830
C      ORDER          - RATE OF CONVERGENCE DESIRED       840

```

```

C      READ(5,102) LEVEL,NRUNS,ORDER          970
C      102  FORMAT(3I2)                      980
C                                             990
C      IF LEVEL = 0 PRINT APPROXIMATE SOLUTION AT NODES 1000
C          = 1 PRINT MAXIMUM ERROR AND MAXIMUM      1010
C              RELATIVE ERROR PROVIDED TRUE SOLUTION
C              IS KNOWN                             1030
C          = 2 ALSO PRINT TRUE SOLUTION AND         1040
C              APPROXIMATE SOLUTION AT THE INTERIOR
C              GRID POINTS                          1050
C                                             1060
C                                             1070
C      IF ORDER = 4 CHOOSE A 4-TH ORDER DIFFERENCE APPR. TO (I) 1080
C          = 6 CHOOSE A 6-TH ORDER DIFFERENCE APPR. TO (I) 1090
C      --STORAGE--IT IS ASSUMED THAT 3 .LE. IXQ,IQY .LE. 7 . 1100
C      IN CASE OF FINER MESH THE DIMENSIONS REQUIRED ARE 1110
C      COMPUTED BY FORMULAS GIVEN IN ARRAY LIST DESCRIPTION. 1120
C      ---- MAIN  VARIABLES OF FFT9 ---- 1130
C                                             1140
C      WORK - WORKING SPACE OF DIMENSION 1150
C      ORDER = 6 7+10*MAX(NGRIDX,NGRIDY)+2*NGRIDX*NGRIDY 1160
C      Z,Y - ARRAYS USED IN THE FOURIER ANALYSIS-SYNTHESIS 1170
C      ORDER = 4 7+10*MAX(NGRIDX,NGRIDY)+NGRIDX*NGRIDY 1180
C              DIMENSION OF 'Z,Y' IS NZD=NYD= MAX(NX,NY) 1190
C      AKX(I+1)-THE I-TH EIGENVALUE OF THE DIAGONAL BLOCK DIVIDED BY 1200
C              THE I-TH EIGENVALUE OF THE OFF DIAGONAL BLOCK 1210
C              SQUARED MINUS 2 1220
C              DIMENSION OF 'AKX' IS NAKXD = MAX(NX,NY)+2 1230
C                                             1240
C                                             1250
C      CORE(K) - VALUE OF APPROXIMATE SOLUTION AT NODE K 1260
C              DIMENSION OF 'CORE' IS NCORED = NX+2+(NX+1)*(NY+1) 1270
C      PTINT( )- VALUES OF G AT HALF LATTICE POINTS 1280
C              DIMENSION IS NPINTD = NX*NY 1290
C      NOTE- ARRAY 'PTINT' IS NOT USED IN CASE ORDER .EQ. 4 1300
C                                             1310
C      GRIDX,GRIDY - GRID COORDINATES 1320
C              DIMENSION OF 'GRIDX,GRIDY' IS 1330
C              NGRDXD = NX+1,NGRDYD = NY+1 1340
C                                             1350
C      SX,SY - LENGTHS OF RECTANGULAR REGION 1360
C                                             1370
C      ED(I+1) - THE I-TH EIGENVALUE OF THE OFF DIAGONAL 1380
C              BLOCK.DIMENSION OF 'ED' IS NEEDD = MAX(NX,NY)+2 1390
C      INDEX - INDEX VECTOR USED IN FFT ANALYSIS 1400
C              DIMENSION IS INDEXD = MAX(NX,NY) 1410
C      SI - SINES VECTOR FOR FFT ANALYSIS AND 1420
C              SYNTHESIS. DIMENSION NSID = MAX(NX,NY) 1430
C                                             1440
C      ---- COMMON VARIABLES OF FFT9 ---- 1450
C                                             1460
C      COMMON /MESH/ 1470
C      NX,NY - NX = 2**IQX,NY = 2**IQY 1480
C      IMIN,IMAX - RANGE OF INTERIOR NODES IN X-DIRECTION 1490
C      JMIN,JMAX - RANGE OF INTERIOR NODES IN Y-DIRECTION 1500
C      INC - INC = IMAX - IMIN + 3 1510
C      IRO - IRO = NX + 3 1520
C      IBCX - IBCX = 1 1530
C      IQX - EXPONENT OF 2 1540
C      IBCY - IBCY = 1 1550
C      IQY - EXPONENT OF 2 1560
C      HX,HY - MESH SIZE 1570
C      HXY2 - HXY2 = (HX/HY)**2 1580
C      PI - PI = 3.14... 1590
C      POTFAC - NORMALIZATION FACTOR = 2/NX 1600
C      COMMON /FDFORM/ 1610
C      DLEFT - DIAGONAL ENTRY OF THE OFFDIAGONAL BLOCK 1620
C      DIAG,OFFD - DIAGONAL AND OFFDIAGONAL ENTRIES OF THE 1630
C              DIAGONAL BLOCK OF THE 9-POINT FORMULA 1640
C      RF,HH,CH,RC1,FACTOR - CONSTANTS 1650
C      COMMON /FFT/ 1660
C      ... - LOCAL VARIABLE 1670
C      ---- COMMON VARIABLES DECLARATION ---- 1680
C                                             1690
C      COMMON /MESH/ NX,NY,IMIN,IMAX,JMIN,JMAX,INC,IRO,IBCX,IQX,IBCY,IQY, 1700
C      IHX,HY,HXY2,PI,POTFAC,SX,SY 1710
C      COMMON /FDFORM/ DLEFT,DIAG,OFFD,RF,HH,CH,RC1,FACTOR 1720

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COMMON /CPDE/ CUXX,CUY,Y,CU          1730
COMMON /FFT/ N2,N4,N3,N7,IP,ISL,L1,IBCJ 1740
COMMON WORK(35040)                   1750
C                                     1760
REAL CVALUS(7)                        1770
INTEGER ORDER                          1780
C                                     1790
C INITIALIZATIONS                     1800
C                                     1810
IBCX=1                                  1820
IBCY=1                                  1830
C                                     1840
C ***** INPUT *****                1850
C DEFINE P.D.E COEFFICIENTS           1860
C                                     1870
CALL PDE (X,Y,CVALUS)                 1880
C                                     1890
CUXX=CVALUS(1)                        1900
CUYY=CVALUS(3)                        1910
CU=CVALUS(6)                          1920
C                                     1930
C DEFINE GRID SPECIFICATIONS           1940
C                                     1950
C READ (5,102) SX,SY,NGRIDX,NGRIDY    1960
NX=NGRIDX-1                            1970
NY=NGRIDY-1                            1980
RNX=NX                                  1990
RNY=NY                                  2000
RALOG2=1./ALOG(2.)                    2010
IQX=ALOG(RNX)*RALOG2                  2020
IQY=ALOG(RNY)*RALOG2                  2030
C                                     2040
C DEFINE OUTPUT CONTROL AND ORDER OF  2050
DESCRITIZATION FORMULA                2060
C                                     2070
C READ (5,103) LEVEL,NRUNS,ORDER      2080
C                                     2090
C OUTPUT THE INPUT DATA               2100
C                                     2110
WRITE (6,104)                          2120
WRITE (6,105)                          2130
WRITE (6,106) CUXX,CUY,Y,CU           2140
WRITE (6,107)                          2150
WRITE (6,108) SX,SY                   2160
WRITE (6,109)                          2170
WRITE (6,110) ORDER                   2180
DO 101 NTIMES=1,NRUNS                 2190
C                                     2200
C ***** DISCRETIZATION *****       2210
C APPROXIMATE THE DIFFERENTIAL EQUATION WITH 9-POINT DIFF.OPER. 2220
C                                     2230
NAKXD=MAX0(NX,NY)+2                   2240
NEDD=NAKXD                             2250
NGRDXD=NX+1                            2260
NGRDYD=NY+1                            2270
C                                     2280
IA1=1                                   2290
IA2=IA1+NAKXD                          2300
IA3=IA2+NEDD                           2310
IA4=IA3+NGRDXD                         2320
IA5=IA4+NGRDYD                         2330
CALL DISCRT (ORDER,WORK(IA1),NAKXD,WORK(IA2),NEDD,WORK(IA3),NGR
1 DXD,WORK(IA4),NGRDYD)                2340
C                                     2350
C                                     2360
C GENERATE RIGHT SIDE OF DIFFERENCE EQUATIONS 2370
C                                     2380
NCORED=NGRDXD*NGRDYD+NAKXD            2390
NPINTD=NCORED                          2400
IF (ORDER.EQ.4) NPINTD=1               2410
IA6=IA5+NCORED                         2420
IA7=IA6+NPINTD                         2430
C                                     2440
CALL RGHTSD (ORDER,WORK(IA5),NCORED,WORK(IA6),NPINTD,WORK(IA3),
1 NGRDXD,WORK(IA4),NGRDYD)            2450
C                                     2460
C ***** EQUATION SOLUTION *****    2470
C                                     2480

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C      2490
C      GENERATE INDECIES AND SINES USED IN THE FOURIER ANALYSIS      2500
C      AND SYNTHESIS      2510
C      2520
C      INEXD=MAX0(NX,NY)      2530
C      NSID=INEXD      2540
C      IA8=IA7+INEXD      2550
C      IA9=IA8+NSID      2560
C      CALL SETF (IBCX,IQX,WORK(IA7),INEXD,WORK(IA8),NSID)      2570
C      2580
C      2590
C      SOLVE THE BLOCK TRIDIAGONAL SYSTEM OF DIFFERENCE EQUATIONS      2600
C      WITH THE FAST FOURIER SERIES METHOD.      2610
C      2620
C      NZD=NAKXD      2630
C      IA10=IA9+NZD      2640
C      NYD=NAKXD      2650
C      CALL EQSOL (WORK(IA5),NCORED,WORK(IA9),NZD,WORK(IA8),NSID,WORK(      2660
1  IA10),NYD,WORK(IA7),INEXD,WORK(IA2),NEDD,WORK(IA1),NAKXD)      2670
C      2680
C      ***** OUTPUT *****      2690
C      2700
C      2710
C      PRINTS THE COMPUTED SOLUTION AND MAX.ERROR,MAX.RELATIVE      2720
C      ERROR IF THE SOLUTION IS KNOWN      2730
C      2740
C      CALL SUMARY (LEVEL,WORK(IA3),NGRDXD,WORK(IA4),NGRDYD,WORK(IA5),      2750
1  NCORED)      2760
C      2770
C      INCREASE EXPONENT OF 2      2780
C      2790
C      IQX=IQX+1      2800
C      IQY=IQY+1      2810
C      NX=2**IQX      2820
C      NY=2**IQY      2830
101 CONTINUE      2840
STOP      2850
C      2860
102 FORMAT (2F10.0,2I3)      2870
103 FORMAT (3I2)      2880
104 FORMAT (1X,27H***** INPUT DATA ***** )      2890
105 FORMAT (1X,37HEQUATION. CUXX*DDXU+CUYY*DDYU+CU*U=R)      2900
106 FORMAT (1X,19HCOEFFICIENTS. CUXX=,F10.3,5HCUY=,F10.3,3HCU=,F10.3)      2910
107 FORMAT (1X,20HBOUNDARY COND. U = G)      2920
108 FORMAT (1X,24HREGION. 0. .LE. X .LE. ,F8.3,3X,15H0. .LE. Y .LE. ,      2930
1F8.3)      2940
109 FORMAT (1X,25H***** SOLUTION ***** )      2950
110 FORMAT (1X,15HDISCRETIZATION.,5X,I5,8H- ORDER ,24HDIFFERENCE APPRO      2960
XIMATION)      2970
C      2980
C      END      2990
C      SUBROUTINE RGHTSD (ORDER,CORE,NCORED,PTINT,NPINTD,GRIDX,NGRDXD,GRI      3000
1DY,NGRDYD)      3010
C      3020
C      COMPUTES RIGHT SIDE OF THE FINITE DIFFERENCE EQUATIONS      3030
C      3040
C      THE ARGUMENTS - ORDER,CORE(NCORED),PTINT(NPINTD),GRIDX(NGRDXD),      3050
C      GRIDY(NGRDYD) - DEFINED IN FFT9 MAIN PROGRAM      3060
C      3070
C      INTEGER ORDER      3080
C      REAL CORE(NCORED),PTINT(NPINTD),GRIDX(NGRDXD),GRIDY(NGRDYD),BVALU      3090
1(4)      3100
C      3110
C      FFT9 COMMON VARIABLES      3120
C      3130
C      COMMON /MESH/ NX,NY,IMIN,IMAX,JMIN,JMAX,INC,IRO,IBCX,IQX,IBCY,IQY,      3140
1HX,HY,HXY2,PI,POTFAC,SX,SY      3150
C      COMMON /FDFORM/ DLEFT,DIAG,OFFD,RF,HH,CH,RC1,FACTOR      3160
C      3170
C      INITIALIZATIONS      3180
C      3190
C      I0=0      3200
C      J0=0      3210
C      Z0=0.0      3220
C      Z1=1.0      3230
C      3240

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C	FUNCTION EVALUATIONS	3250
C		3260
	DO 102 J=J0,NY	3270
	W=GRIDY(J+1)	3280
	L=IRO+INC*J	3290
	DO 101 I=I0,NX	3300
	K=L+I	3310
	X=GRIDX(I+1)	3320
	CORE(K)=PDERGH(X,W)	3330
101	CONTINUE	3340
102	CONTINUE	3350
	K=IRO	3360
C		3370
C	CORNER INDECIES	3380
C		3390
	NR=IRO+NX	3400
	NU=IRO+NY*INC	3410
	NRU=NU+NX	3420
C		3430
	IF (ORDER.EQ.4) GO TO 108	3440
C		3450
C	INITIALIZATIONS	3460
C		3470
	DO 103 I=I0,NX	3480
	CORE(I+1)=CORE(K)	3490
	K=K+1	3500
103	CONTINUE	3510
C		3520
C	EVALUATE RIGHT SIDE AT THE EXTRA POINTS NEEDED BY SIX	3530
C	ORDER FORMULA	3540
C		3550
	YMIDL=-HY*.5	3560
	DO 105 J=1,NY	3570
	L=NX*(J-1)	3580
	YMIDL=YMIDL+HY	3590
	XMIDL=-HX*.5	3600
	DO 104 I=1,NX	3610
	XMIDL=XMIDL+HX	3620
	K=I+L	3630
	PTINT(K)=PDERGH(XMIDL,YMIDL)	3640
104	CONTINUE	3650
105	CONTINUE	3660
C		3670
C	COMPUTE THE RIGHT SIDE OF SIX ORDER DIFFERENCE OPERATOR	3680
C		3690
	LL=IRO	3700
	DO 107 J=JMIN,JMAX	3710
	LL=LL+INC	3720
	L=LL	3730
	CNTRLF=CORE(L)	3740
	LDOWN=L-INC	3750
	DOWNLF=CORE(LDOWN)	3760
	DO 106 I=IMIN,IMAX	3770
	LUP=L+INC	3780
	LUP2=LUP+2	3790
	LUP1=LUP+1	3800
	IP1=I+1	3810
	IP2=I+2	3820
	LP2=L+2	3830
	K=L+1	3840
	TEMP=DOWNLF+CORE(LUP)+CORE(LUP2)+CORE(IP2)+4.*(CNTRLF+CORE(L	3850
1	UP1)+CORE(LP2)+CORE(IP1))+148.*CORE(K)	3860
C		3870
	CNTRLF=CORE(K)	3880
	DOWNLF=CORE(IP1)	3890
	CORE(I+1)=CORE(K)	3900
	IDWN1=I+NX*(J-1)	3910
	IDWN2=IDWN1+1	3920
C		3930
	IUP1=IDWN1+NX	3940
	IUP2=IUP1+1	3950
C		3960
	CORE(K)=FACTOR*(TEMP+48.*(PTINT(IDWN1)+PTINT(IDWN2)+PTINT(IU	3970
1	P1)+PTINT(IUP2)))	3980
	L=L+1	3990
106	CONTINUE	4000


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          LP1=L+1
          CORE(NX+1)=CORE(LP1)
107 CONTINUE
C
      GO TO 112
108 CONTINUE
C
C   INITIALIZATION
C
      DO 109 I=IMIN,IMAX
          K=K+1
          CORE(I)=CORE(K)
109 CONTINUE
C
C   GENERATE RIGHT SIDE OF DIFFERENCE EQUATIONS
C
      L=IRO
      DO 111 J=JMIN,JMAX
          L=L+INC
          K=L
          XCENR=CORE(I)
          KRIGHT=K+1
          DO 110 I=IMIN,IMAX
              K=KRIGHT
              KRIGHT=K+1
              KUP=K+INC
              XLEFT=XCENR
              XRIGHT=CORE(KRIGHT)
              XVERT=CORE(I)+CORE(KUP)
              XCENR=CORE(K)
              CORE(I)=XCENR
              CORE(K)=FACTOR*(XVERT+HH*(XLEFT+XRIGHT)+CH*XCENR)
110 CONTINUE
111 CONTINUE
C
112 CONTINUE
C
C   BOUNDARY VALUES AT THE CORNERS
C
      CORE(IRO)=BCOND(3,0.,0.,BVALUS)
      CORE(NR)=BCOND(2,SX,0.,BVALUS)
      CORE(NU)=BCOND(4,0.,SY,BVALUS)
      CORE(NRU)=BCOND(1,SX,SY,BVALUS)
C
C   ENFORCE DIRICHLET BOUNDARY CONDITIONS
C
      KLEFT=IRO
      KLFTUP=IRO+INC
      MRIGHT=NR
      MRGTUP=MRIGHT+INC
      CORE(KLFTUP)=BCOND(3,0.,HY,BVALUS)
      CORE(MRGTUP)=BCOND(1,SX,HY,BVALUS)
      DO 113 J=JMIN,JMAX
          W=GRIDY(J+2)
          KLEFTD=KLEFT
          KLEFT=KLFTUP
          KLFTUP=KLFTUP+INC
          K=KLEFT+1
          MRGTD=MRIGHT
          MRIGHT=MRGTUP
          MRGTUP=MRGTUP+INC
          M=MRIGHT-1
          CORE(KLFTUP)=BCOND(3,0.,W,BVALUS)
          CORE(MRGTUP)=BCOND(1,SX,W,BVALUS)
          CORE(K)=CORE(K)-POTFAC*(CORE(KLEFTD)+CORE(KLFTUP)+OFFD*CORE(KLE
1          FT))
          CORE(M)=CORE(M)-POTFAC*(CORE(MRGTD)+CORE(MRGTUP)+OFFD*CORE(MRIG
1          HT))
113 CONTINUE
      KDOWN=IRO
      KRGTD=IRO+1
      MUP=NU
      MRGTUP=NU+1
      CORE(KRGTD)=BCOND(2,HX,0.,BVALUS)
      CORE(MRGTUP)=BCOND(4,HX,SY,BVALUS)

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C          5530
      DO 104 K=IMIN,IMAX          5540
        A=AKX(K+1)                5550
        L=IRO+K                    5560
        M=JUMP                      5570
104 CALL CRED (IBCY,L,M,A,IQY-1,CORE,NCORED) 5580
C          5590
C      FOURIER SYNTHESIS ON EVEN LINES 5600
C          5610
      K=IRO+INC*J1                5620
      DO 105 J=J1,J2,2            5630
        CALL FETCHX (K,Z,NZD,CORE,NCORED) 5640
        CALL FOUR (IBCX,IQX,SI,NSID,Z,NZD,Y,NYD,INDEX,INDEXD) 5650
        CALL STOREX (K,CORE,NCORED,Y,NYD) 5660
105 K=K+JUMP                      5670
C          5680
C      MODIFICATION OF ODD LINE VECTORS 5690
C          5700
      J2=NY-1                      5710
      F1=1./POTFAC                 5720
      CALL ODDRD (F1,J2,CORE,NCORED) 5730
C          5740
C      SOLUTION FOR ODD LINES BY CYCLIC REDUCTION 5750
C          5760
      A=-DIAG/OFFD                 5770
      DO 106 J=1,J2,2              5780
        L=IRO+INC*J                5790
106 CALL CRED (IBCX,L,1,A,IQX,CORE,NCORED) 5800
C          5810
      RETURN                        5820
C          5830
      END                            5840
      SUBROUTINE DISCRT (ORDER,AKX,NAKXD,ED,NEDD,GRIDX,NGRDXD,GRIDY,NGRD
1YD)                                5850
C          5860
C          5870
C      SETS CONSTANTS,CALCULATE GRID SPECIFICATIONS,9-POINT 5880
C      DIFFERENCE FORMULA AND THE EIGENVALUES OF THE DIAGONAL 5890
C      AND OFF DIAGONAL BLOCKS OF THE FINITE DIFFERENCE EQUATIONS 5900
C          5910
C      THE ARGUMENTS-ORDER,AKX(NAKXD),ED(NEDD),GRIDX(NGRDXD), 5920
C      GRIDY(NGRDYD) ARE DEFINED IN FFT9 5930
C          5940
C      INTEGER ORDER                5950
C      REAL AKX(NAKXD),ED(NEDD),GRIDX(NGRDXD),GRIDY(NGRDYD) 5960
C          5970
C      FFT9 COMMON VARIABLES        5980
C          5990
      COMMON /MESH/ NX,NY,IMIN,IMAX,JMIN,JMAX,INC,IRO,IBCX,IQX,IBCY,IQY,
1HX,HY,HXY2,PI,POTFAC,SX,SY        6000
      COMMON /FDFORM/ DLEFT,DIAG,OFFD,RF,HH,CH,RC1,FACTOR 6010
      COMMON /CPDE/ C1,C2,C3        6020
C          6030
C          6040
C      GENERATE CONSTANTS          6050
C          6060
      PI=3.14159265358979          6070
      NX=2**IQX                     6080
      NY=2**IQY                     6090
      REV=1./FLOAT(NX)              6100
      REVY=1./FLOAT(NY)             6110
      POTFAC=2.*REV                 6120
      HX=SX*REV                     6130
      HY=SY*REVY                    6140
      RF=PI*REV                     6150
C          6160
C      GENERATE GRID SPECIFICATIONS 6170
C          6180
      I0=0                          6190
      DO 101 I=I0,NX                 6200
        GRIDX(I+1)=FLOAT(I)*HX     6210
101 CONTINUE                       6220
      DO 102 J=I0,NY                 6230
        GRIDY(J+1)=FLOAT(J)*HY     6240
102 CONTINUE                       6250
      IMIN=1                        6260
      IMAX=NX-1                     6270
      JMIN=1                         6280

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        JMAX=NY-1
        INC=IMAX-IMIN+3
        IRO=2+INC
        HXY2=(HX/HY)**2
C
C   INITIALIZATIONS
C
        OFFD=4.
        DLEFT=4.
        DIAC=-20.
        FACTOR=POTFAC*HX**2/60.
C
        IF (ORDER.EQ.6) GO TO 103
C
C   GENERATE COEFFICIENTS OF 9-POINT FINITE DIFFERENCE STENSIL
C   OF FOURTH ORDER
C
        RC1=1./C1
        C21=RC1*C2
        SIGMA=RC1*C3
        RGRID=HY/HX
        HX2=HX**2
        SIGH2=SIGMA*HX2
        SIGH12=SIGH2/12.0
        RGRSQ=RGRID**2
        QUOT=RGRSQ/C21
        RR=1.-SIGH12
        QQ=C21*(1.-SIGH12*QUOT)/RR
        DIV=RGRSQ+QQ
        OFFD=12.*QQ*QUOT/DIV-2.0
        DLEFT=12.*C21/DIV-2.0
        DIAC=(OFFD+2.)*RR*SIGH2-2.*(OFFD+DLEFT)-4.
        HH=QQ/C21
        CH=(12.*RR-2.)*HH-2.
        FACTOR=RC1*POTFAC*RGRSQ*HX2/DIV
103 CONTINUE
C
C   CALCULATE EIGENVALUES OF THE OFF-DIAGONAL BLOCKS AND
C   DIAGONAL BLOCKS
C
        DO 104 I=IMIN,IMAX
            II=I+1
            DFLI=FLOAT(I)
            TWOCOS=2.0*COS(RF*DFLI)
            ED(II)=(DLEFT+TWOCOS)**2
            RATIO=((DIAC+TWOCOS*OFFD)**2)/ED(II)
            AKX(II)=-2.0+RATIO
104 CONTINUE
        RETURN
C
        END
        SUBROUTINE EVENRD (J1,J2,CORE,NCORED)
C
C   MODIFIES THE RIGHT SIDE ON EVEN LINE VECTORS WHERE J1 IS
C   THE FIRST AND J2 THE LAST EVEN VECTOR.
C   THE RIGHT SIDE AND ITS MODIFICATION ARE STORED IN ARRAY CORE
C
C   THE ACTUAL ARGUMENT - CORE(NCORED) IS DEFINED IN FFT9
C
        REAL CORE(NCORED)
C
C   FFT9 COMMON VARIABLES
C
        COMMON /MESH/ NX,NY,IMIN,IMAX,JMIN,JMAX,INC,IRO,IBCX,IQX,IBCY,IQY,
        1HX,HY,HXY2,PI,POTFAC,SX,SY
        COMMON /FDFORM/ DLEFT,DIAC,OFFD,RF,HH,CH,RC1,FACTOR
C
C   INITIALIZATIONS
C
        I1=IMIN
        I2=IMAX-1
        L=IRO
        JUMP=INC+INC
C
        DO 102 J=J1,J2,2
            L=L+JUMP

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      K=L
      KRIGHT=K+1
      KRDOWN=KRRIGHT-INC
      KRGTUP=KRRIGHT+INC
      X2=0.
      X3=CORE (KRDOWN)+CORE (KRGTUP)
      Y2=0.
      Y3=CORE (KRIGHT)
C
      DO 101 I=I1,I2
      K=K+1
      KRIGHT=KRRIGHT+1
      KRDOWN=KRDOWN+1
      KRGTUP=KRGTUP+1
      X1=X2
      X2=X3
      X3=CORE (KRDOWN)+CORE (KRGTUP)
      Y1=Y2
      Y2=Y3
      Y3=CORE (KRIGHT)
      CORE1=X1+X3+DLEFT*X2-OFFD*(Y1+Y3)-DIAG*Y2
      CORE(K)=CORE1
101  CONTINUE
      K=K+1
      CORE1=X2+DLEFT*X3-OFFD*Y2-DIAG*Y3
      CORE(K)=CORE1
102  CONTINUE
      RETURN
C
      END
      SUBROUTINE ODDRD (F1,J2,CORE,NCORED)
C
C   MODIFIES THE RIGHT SIDE ON ODD-LINE VECTORS WHERE J2 IS THE LAST
C   ODD LINE.THE RIGHT SIDE AND ITS MODIFICATION IS STORED IN
C   ARRAY CORE(NCORED).
C
C   THE ARGUMENT F1 IS A MULTIPLICATION FACTOR
C
C   REAL CORE(NCORED)
C
C   FFT9 COMMON VARIABLES
C
      COMMON /MESH/ NX,NY,IMIN,IMAX,JMIN,JMAX,INC,IRO,IBCX,IQX,IBCY,IQY,
      LHX,HY,HXY2,PI,POTFAC,SX,SY
      COMMON /FDFORM/ DLEFT,DIAG,OFFD,RF,HH,CH,RC1,FACTOR
C
      DENOM=1./OFFD
      DLFT=DLEFT*DENOM
      CENTER=F1*DENOM
      I2=IMAX-1
      L=IRO-INC
      JUMP=INC+INC
C
      DO 102 J=1,J2,2
      L=L+JUMP
      K=L
      KRIGHT=K+1
      KRDOWN=KRRIGHT-INC
      KRGTUP=KRRIGHT+INC
      X2=0.
      X3=CORE (KRDOWN)+CORE (KRGTUP)
      IF (J.EQ.1) X3=CORE (KRGTUP)
      IF (J.EQ.J2) X3=CORE (KRDOWN)
      DO 101 I=IMIN,I2
      K=K+1
      KRIGHT=KRRIGHT+1
      KRDOWN=KRDOWN+1
      KRGTUP=KRGTUP+1
      X1=X2
      X2=X3
      X3=CORE (KRDOWN)+CORE (KRGTUP)
      IF (J.EQ.1) X3=CORE (KRGTUP)
      IF (J.EQ.J2) X3=CORE (KRDOWN)
      CORE1=CENTER*CORE(K)-DENOM*(X1+X3)-DLFT*X2
      CORE(K)=CORE1
101  CONTINUE

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C		7810
	K=K+1	7820
	CORE1=CENTER*CORE(K)-DENOM*X2-DLFT*X3	7830
	CORE(K)=CORE1	7840
102	CONTINUE	7850
	RETURN	7860
C		7870
	END	7880
	SUBROUTINE CRED (IBC,L,M,A,IP1,CORE,NCORED)	7890
C		7900
C	SOLVES TRIANGULAR SYSTEMS BY RECURSIVE CYCLIC REDUCTION	7910
C	SEE REFERENCE @3	7920
C		7930
C	IBC = 1,IP1-EXPONENT OF 2,A-DIAGONAL ELEMENT,L,M-MESH DEPENDENT	7940
C	CONSTANTS USED TO RECOVER RIGHT SIDE FROM ARRAY CORE	7950
C		7960
	REAL CORE(NCORED)	7970
	COMMON /MESH/ NX,NY,IMIN,IMAX,JMIN,JMAX,INC,IRO,IBCX,IQX,IBCY,IQY,	7980
	LHX,HY,HXY2,PI,POTFAC,SX,SY	7990
	DIMENSION BB(11)	8000
C		8010
	IP=IP1	8020
	N2=M	8030
	BB(1)=A	8040
	B=A	8050
	N4=0	8060
	N=2**IP	8070
	K=L+N*M	8080
	IP=IP-1	8090
C		8100
	DO 104 N1=1,IP	8110
	N4=N4+1	8120
	N3=N2	8130
	N2=N2+N2	8140
	J1=N2+L	8150
	J3=K-N3	8160
	J2=K-N2	8170
	IF (J1.GT.J2) GO TO 102	8180
C		8190
	DO 101 J=J1,J2,N2	8200
	JMN3=J-N3	8210
	JPN3=J+N3	8220
	CORE(J)=B*CORE(J)+CORE(JMN3)+CORE(JPN3)	8230
101	CONTINUE	8240
102	B=B*B-2.000	8250
	BB(N1+1)=B	8260
	IF (B.LE.1.0E14) GO TO 104	8270
C		8280
C	SHORT CUT AND SOLVE BY DIVISION IF B	8290
C	LARGER THAN DESIRED ACCURACY	8300
C		8310
	IF (J1.GT.J2) GO TO 105	8320
C		8330
	DO 103 J=J1,J2,N2	8340
103	CORE(J)=-CORE(J)/B	8350
C		8360
	IF (IBC.EQ.1) GO TO 105	8370
104	CONTINUE	8380
	I=L+N*M/2	8390
	CORE(I)=-CORE(I)/B	8400
C		8410
105	DO 108 NN=1,N4	8420
	N1=N4-NN	8430
	B=BB(N1+1)	8440
	J2=K-N3	8450
	J1=L+N3	8460
	J1PN3=J1+N3	8470
	CORE(J1)=(CORE(J1PN3)-CORE(J1))/B	8480
	J2MN3=J2-N3	8490
	CORE(J2)=(CORE(J2MN3)-CORE(J2))/B	8500
	J1=J1+N2	8510
	J2=J2-N2	8520
	IF (J1.GT.J2) GO TO 107	8530
C		8540
	DO 106 J=J1,J2,N2	8550
	JMN3=J-N3	8560

```

          JPN3=J+N3
          CORE(J)=(CORE(JMN3)+CORE(JPN3)-CORE(J))/B
106  CONTINUE
107  N2=N3
108  N3=N3/2
      RETURN
C
      END
      SUBROUTINE STOREX (K,CORE,NCORED,Y,NYD)
C
C   TRANSFERS DATA FROM ARRAY Y TO ARRAY CORE
C   AFTER FOURIER ANALYSIS
C
C   THE ARGUMENTS - K IS THE NODE,CORE(NCORED),Y(NYD) DEFINED IN FFT9
C
      REAL CORE(NCORED),Y(NYD)
      COMMON /MESH/ NX,NY,IMIN,IMAX,JMIN,JMAX,INC,IRO,IBCX,IQX,IBCY,IQY,
1HX,HY,HXY2,PI,POTFAC,SX,SY
      DO 101 I=IMIN,IMAX
          KPI=K+I
          IP1=I+1
          CORE(KPI)=Y(IP1)
101  CONTINUE
      RETURN
C
      END
      SUBROUTINE FETCHX (K,Z,NZD,CORE,NCORED)
C
C   TRANSFERS DATA FROM ARRAY CORE TO ARRAY Z PRIOR TO
C   FOURIER ANALYSIS
C
C   THE ARGUMENTS - K NODE,Z(NZD),CORE(NCORED) ARE DEFINED IN FFT9
C
      REAL Z(NZD),CORE(NCORED)
      COMMON /MESH/ NX,NY,IMIN,IMAX,JMIN,JMAX,INC,IRO,IBCX,IQX,IBCY,IQY,
1HX,HY,HXY2,PI,POTFAC,SX,SY
      DO 101 I=IMIN,IMAX
          KPI=K+I
          IP1=I+1
          Z(IP1)=CORE(KPI)
101  CONTINUE
      RETURN
C
      END
      SUBROUTINE KFOLD (INDEX,INDEXD,SI,NSID,Y,NYD,Z,NZD)
C
C   EVALUATES THE SUMMATIONS AND DOES ALL THE MULTIPLICATIONS
C   BY A RECURSIVE TECHNIQUE (SEE REFERENCE @3 )
C
C   THE ARGUMENTS--INDEX(INDEXD),SI(NSID),Y(NYD),Z(NZD) DEFINED IN FFT9
C
      REAL SI(NSID),Y(NYD),Z(NZD)
      INTEGER INDEX(INDEXD)
C
C   FFT9 COMMON VARIABLES
C
      COMMON /FFT/ N2,N4,N3,N7,IP,ISL,L1,IBCJ
C
      JS1=N2
      I=1
      J5=ISL+N2
      IS1=ISL
      ICL=L1
      JS1=JS1/2
C
C   GO TO FIRST TIME IS LAST TIME
C
      IF (JS1.EQ.1) GO TO 106
      SN=SI(1)
      IS1=IS1+JS1
      ICL=ICL+JS1
      J3=IS1+JS1
C
101  ISO=IS1-JS1
      ICO=ICL-JS1
      ODD1=SN*(Z(ICL)-Z(IS1))

```

```

        ODD2=SN*(Z(IC1)+Z(IS1))          9330
        Z(IC1)=Z(ICO)-ODD1              9340
        Z(ICO)=Z(ICO)+ODD1              9350
        Z(IS1)=-Z(ISO)+ODD2             9360
        Z(ISO)=Z(ISO)+ODD2              9370
        IS1=IS1+1                       9380
        IC1=IC1+1                       9390
        IF (IS1.NE.J3) GO TO 101         9400
C                                          9410
        I=I+1                           9420
102 IS1=ISL                             9430
        IC1=L1                           9440
        JS1=JS1/2                       9450
C                                          9460
C    GO TO LAST TIME WITH K IN PAIRS    9470
C                                          9480
        IF (JS1.EQ.1) GO TO 107         9490
C                                          9500
C    TAKE K IN PAIRS INTERCHANGING SN AND CS 9510
C                                          9520
103 SN=SI(I)                            9530
        I=I+1                           9540
        CS=SI(I)                        9550
        IS1=IS1+JS1                     9560
        IC1=IC1+JS1                     9570
        J3=IS1+JS1                      9580
C                                          9590
104 ISO=IS1-JS1                         9600
        ICO=IC1-JS1                     9610
        ODD1=CS*Z(IC1)-SN*Z(IS1)        9620
        ODD2=SN*Z(IC1)+CS*Z(IS1)        9630
        Z(IC1)=Z(ICO)-ODD1              9640
        Z(ICO)=Z(ICO)+ODD1              9650
        Z(IS1)=-Z(ISO)+ODD2             9660
        Z(ISO)=Z(ISO)+ODD2              9670
        IS1=IS1+1                       9680
        IC1=IC1+1                       9690
        IF (IS1.NE.J3) GO TO 104         9700
C                                          9710
        IS1=IS1+JS1                     9720
        IC1=IC1+JS1                     9730
        J3=IS1+JS1                     9740
C                                          9750
105 ISO=IS1-JS1                         9760
        ICO=IC1-JS1                     9770
        ODD1=SN*Z(IC1)-CS*Z(IS1)        9780
        ODD2=CS*Z(IC1)+SN*Z(IS1)        9790
        Z(IC1)=Z(ICO)-ODD1              9800
        Z(ICO)=Z(ICO)+ODD1              9810
        Z(IS1)=-Z(ISO)+ODD2             9820
        Z(ISO)=Z(ISO)+ODD2              9830
        IS1=IS1+1                       9840
        IC1=IC1+1                       9850
        IF (IS1.NE.J3) GO TO 105         9860
        I=I+1                           9870
        IF (IS1.EQ.J5) GO TO 102         9880
        GO TO 103                       9890
C                                          9900
C    LAST TIME IS FIRST TIME           9910
C                                          9920
106 K1=INDEX(I)                         9930
        SN=SI(I)                        9940
        ISO=IS1                          9950
        IS1=IS1+JS1                     9960
        ICO=IC1                          9970
        IC1=IC1+JS1                     9980
        ODD1=SN*(Z(IC1)-Z(IS1))         9990
        Y(K1+1)=Z(ICO)+ODD1            10000
        NDUM=N3-K1+1                    10010
        Y(NDUM)=Z(ICO)-ODD1             10020
        RETURN                           10030
C                                          10040
C    LAST TIME TAKING K IN PAIRS       10050
C                                          10060
107 K1=INDEX(I)                         10070
        SN=SI(I)                        10080

```



```

I=I+1 10090
CS=SI (I) 10100
ISO=IS1 10110
IS1=IS1+JS1 10120
ICO=IC1 10130
IC1=IC1+JS1 10140
ODD1=CS*Z (IC1)-SN*Z (IS1) 10150
Y (K1+1)=Z (ICO)+ODD1 10160
NDUM=N3-K1+1 10170
Y (NDUM)=Z (ICO)-ODD1 10180
C 10190
IS1=IS1+1 10200
IC1=IC1+1 10210
K1=INDEX (I) 10220
ISO=IS1 10230
IS1=IS1+JS1 10240
ICO=IC1 10250
IC1=IC1+JS1 10260
ODD1=SN*Z (IC1)-CS*Z (IS1) 10270
Y (K1+1)=Z (ICO)+ODD1 10280
NDUM=N3-K1+1 10290
Y (NDUM)=Z (ICO)-ODD1 10300
C 10310
IS1=IS1+1 10320
IC1=IC1+1 10330
I=I+1 10340
IF (IS1.NE.J5) GO TO 107 10350
RETURN 10360
C 10370
END 10380
SUBROUTINE FOUR (IBC1,IQ1,SI,NSID,Z,NZD,Y,NYD,INDEX,INDEXD) 10390
C 10400
C PERFORMS A SINE ANALYSIS OR SYNTHESIS ON THE INPUT 10410
C ARRAY Z (I) , I = 2,N AND PUTS THE RESULTS IN THE ARRAY Y 10420
C 10430
C THE ARGUMENTS - ALL - DEFINED IN FFT9 MAIN PROGRAM 10440
C 10450
C REAL SI (NSID) , Z (NZD) , Y (NYD) 10460
C INTEGER INDEX (INDEXD) 10470
C 10480
C FFT9 COMMON VARIABLES 10490
C 10500
C COMMON /FFT/ N2,N4,N3,N7,IP,IS1,L1,IBCK 10510
C 10520
C IBC=IBC1 10530
C IQ=IQ1 10540
C A5=SI (1) 10550
C N4=2**IQ 10560
C N3=N4 10570
C N5=N3/4 10580
C N7=N3/2 10590
C N11=3*N7 10600
C N31=N3+1 10610
C Z (N31)=0.000 10620
C Z (1)=0.000 10630
C N2=N3 10640
C 10650
C DO 101 I=2,IQ 10660
C 10670
C CALL TFOLD (1,1,Z,NZD) 10680
101 N2=N2/2 10690
Y (N7+1)=Z (2) 10700
JF=N5 10710
C 10720
C DO 102 J=2,IQ 10730
C L1=N2+1 10740
C 10750
C CALL ZERO (1,Z,NZD) 10760
C IS1=1 10770
C 10780
C CALL KFOLD (INDEX,INDEXD,SI,NSID,Y,NYD,Z,NZD) 10790
C I1=3*JF+1 10800
C I2=4*JF 10810
C I3=I1+(N2/2-1)*I2 10820
C 10830
C CALL NEG (I1,I3,I2,Y,NYD) 10840

```

```

      N2=N2+N2
102 JF=JF/2
      RETURN
C
      END
      SUBROUTINE NEG (I1,I3,I2,Y,NYD)
C
C   SETS THE Y(K) K=I1,I3,I2 EQUAL TO -Y(K)
C
C   THE ARGUMENT Y(NYD) IS DEFINED IN FFT9
C
      REAL Y(NYD)
      DO 101 K=I1,I3,I2
101 Y(K)=-Y(K)
      RETURN
C
      END
      SUBROUTINE ZERO (L,Z,NZD)
C
C   SETS THE ELEMENTS OF Z(LM)=0 , LM=L+I-1, FOR I=1,N2
C
      REAL Z(NZD)
      COMMON /FFT/ N2,N4,N3,N7,IP,ISL,L1,IBCM
      DO 101 I=1,N2
          LDUM=L+I-1
          Z(LDUM)=0.0
101 CONTINUE
      RETURN
C
      END
      SUBROUTINE TFOLD (IS,L,Z,NZD)
C
C   FOLDS THE INPUT ARRAY Z(NZD) DEFINED IN FFT9
C
      REAL Z(NZD)
      COMMON /FFT/ N2,N4,N3,N7,IP,ISL,L1,IBCN
      IH2=N2/2-1
      DO 101 I=IS,IH2
          I1=I+L
          I2=N2-I+L
          A=Z(I1)
          Z(I1)=A-Z(I2)
101 Z(I2)=A+Z(I2)
      RETURN
C
      END
      SUBROUTINE SETF (IBC1,IQ1,INDEX,INDEXD,SI,NSID)
C
C   INITIALIZES ARRAYS SI ,INDEX FOR THE USE
C   BY FOUR( SEE REFERENCE @3 )
C
C   THE ARGUMENTS - ALL - DEFINED IN FFT9 MAIN PROGRAM
C
      REAL SI(NSID)
      INTEGER INDEX(INDEXD)
      COMMON /FFT/ N2,N4,N3,N7,IP,ISL,L1,IBCO
      DPI=3.14159265358979
      IBC=IBC1
      IQ=IQ1
      N3=2**IQ
      N7=N3/2
      N5=N3/4
      I=1
      INDEX(I)=N5
      SI(I)=1.0/SQRT(2.0)
      K=1
      I=I+1
101 IL=I
      IF (I.EQ.N7) GO TO 103
102 K1=INDEX(K)/2
      INDEX(I)=K1
      RK1=K1
      RN3=N3
      SI(I)=SIN(DPI*RK1/RN3)
      K1=N7-K1
      I=I+1

```

```

10850
10860
10870
10880
10890
10900
10910
10920
10930
10940
10950
10960
10970
10980
10990
11000
11010
11020
11030
11040
11050
11060
11070
11080
11090
11100
11110
11120
11130
11140
11150
11160
11170
11180
11190
11200
11210
11220
11230
11240
11250
11260
11270
11280
11290
11300
11310
11320
11330
11340
11350
11360
11370
11380
11390
11400
11410
11420
11430
11440
11450
11460
11470
11480
11490
11500
11510
11520
11530
11540
11550
11560
11570
11580
11590
11600

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INDEX(I)=K1 11610
RK1=K1 11620
RN3=N3 11630
SI(I)=SIN(DPI*RK1/RN3) 11640
K=K+1 11650
I=I+1 11660
IF (K.NE.II) GO TO 102 11670
GO TO 101 11680
103 RETURN 11690
C 11700
END 11710
C ***** OUTPUT MODULU ***** 11720
C 11730
SUBROUTINE SUMARY (LEVEL,GRIDX,NGRDxD,GRIDY,NGRDYD,CORE,NCORED) 11740
C 11750
PRINTS THE COMPUTED SOLUTION AND MAX. ERROR,MAX. 11760
RELATIVE ERROR IF THE SOLUTION IS KNOWN 11770
THE ARGUMENTS - ALL - DEFINED IN FFT9 MAIN PROGRAM 11780
C 11790
COMMON VARIABLES OF FFT9 11800
C 11810
REAL GRIDX(NGRDxD),GRIDY(NGRDYD),CORE(NCORED) 11820
COMMON /MESH/ NX,NY,IMIN,IMAX,JMIN,JMAX,INC,IRO,IBCX,IQX,IBCY,IQY, 11830
1HX,HY,HXY2,PI,POTFAC,SX,SY 11840
COMMON /CPDE/ CUXX,CUYX,CU 11850
C 11860
C 11870
MSLINX=NX+1 11880
MSLINY=NY+1 11890
C 11900
WRITE (6,108) MSLINX,MSLINY 11910
IF (LEVEL.GT.0) GO TO 103 11920
WRITE (6,109) 11930
DO 102 I=IMIN,IMAX 11940
DO 101 J=JMIN,JMAX 11950
K=IRO+I+INC*J 11960
WRITE (6,110) GRIDX(I+1),GRIDY(J+1),CORE(K) 11970
101 CONTINUE 11980
102 CONTINUE 11990
C 12000
RETURN 12010
C 12020
C COMPUTE MAX.ABSOLUTE AND MAX.RELATIVE ERROR 12030
C 12040
C 12050
103 CONTINUE 12060
IF (LEVEL.EQ.2) WRITE (6,111) 12070
ERRMAX=0. 12080
RELMAX=0. 12090
DO 107 I=IMIN,IMAX 12100
DO 106 J=JMIN,JMAX 12110
K=IRO+I+INC*J 12120
EXCT=TRUE(GRIDX(I+1),GRIDY(J+1)) 12130
ERROR=ABS(CORE(K)-EXCT) 12140
IF (ERROR.LE.ERRMAX) GO TO 104 12150
ERRMAX=ERROR 12160
XMAX=GRIDX(I+1) 12170
YMAX=GRIDY(J+1) 12180
104 CONTINUE 12190
IF (ABS(EXCT).LT.1.E-14) EXCT=1.E-14 12200
RELERR=ERROR/ABS(EXCT) 12210
IF (RELERR.LE.RELMAX) GO TO 105 12220
RELMAX=RELERR 12230
XREL=GRIDX(I+1) 12240
YREL=GRIDY(J+1) 12250
105 CONTINUE 12260
C 12270
C PRINT SOLUTION ,APPROXIMATE SOLUTION ,MAX.ERROR,REL.ERROR 12280
C AT EACH POINT 12290
C 12300
IF (LEVEL.EQ.2) WRITE (6,112) GRIDX(I+1),GRIDY(J+1),EXCT,COR 12310
E(K),ERROR,RELERR 12320
106 CONTINUE 12330
107 CONTINUE 12340
WRITE (6,113) ERRMAX,XMAX,YMAX,RELMAX,XREL,YREL 12350
RETURN 12360

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ALGORITHM 544

L2A and L2B, Weighted Least Squares Solutions by Modified Gram-Schmidt with Iterative Refinement [F4]

ROY H. WAMPLER
National Bureau of Standards

Key Words and Phrases: covariance matrix, curve fitting, iterative refinement, least squares solution, linear constraints, overdetermined system of equations, regression, underdetermined system of equations
CR Categories: 5.14, 5.5
Language: PFORT, a portable subset of ANSI Fortran

DESCRIPTION

The two algorithms given here are a complement to [1] where the types of problems which can be solved by L2A and L2B, together with the method of solution and details of the calling sequences, are described.

REFERENCES

1. WAMPLER, R.H. Solutions to weighted least squares problems by modified Gram-Schmidt with iterative refinement. *ACM Trans. Math. Software* 5, 4(Dec. 1979), 457-465.

ALGORITHM

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C MAIN PROGRAM TO CALL SUBROUTINES L2A AND L2B FOR SOLVING LINEAR      DRV  10
C LEAST SQUARES PROBLEMS.                                             DRV  20
C                                                                       DRV  30
C VERSION OF NOVEMBER 2, 1978.                                        DRV  40
C                                                                       DRV  50
C WRITTEN BY ROY H. WAMPLER, STATISTICAL ENGINEERING                 DRV  60
C LABORATORY, NATIONAL BUREAU OF STANDARDS,                          DRV  70
C WASHINGTON, D. C. 20234.                                           DRV  80
C                                                                       DRV  90
C THE MAIN PROGRAM READS AND PRINTS THE INPUT DATA, CALLS EITHER   DRV 100
C SUBROUTINE L2A OR L2B, COMPUTES CERTAIN QUANTITIES DERIVED FROM   DRV 110
C OUTPUT OF THE SUBROUTINE, AND PRINTS COMPUTED RESULTS.           DRV 120
C                                                                       DRV 130
C SEE COMMENTS AT THE BEGINNING OF SUBROUTINES L2A AND L2B FOR A    DRV 140
C DESCRIPTION OF THE INPUT, OUTPUT AND INTERNAL VARIABLES WHICH APPEAR DRV 150
C IN THE CALLS TO THOSE SUBROUTINES.                                  DRV 160
C                                                                       DRV 170
C MODE IS A PARAMETER WHICH CONTROLS WHETHER SUBROUTINE L2A OR     DRV 180
C SUBROUTINE L2B SHALL BE CALLED BY THIS MAIN PROGRAM. THE TWO     DRV 190
C SUBROUTINES WILL FURNISH THE SAME SOLUTIONS WHENEVER THE COMPUTED DRV 200
C RANK OF THE SYSTEM OF M EQUATIONS IN N UNKNOWNNS EQUALS N AND M.G.E.N. DRV 210
C IN CASES WHERE THE COMPUTED RANK N1 IS LESS THAN N, THE USER     DRV 220
C SPECIFIES THE TYPE OF SOLUTION TO BE COMPUTED ACCORDING TO WHETHER DRV 230
C MODE = 1 OR MODE = 2.                                             DRV 240
C MATRIX A IS THE GIVEN MATRIX OF A SYSTEM OF M LINEAR EQUATIONS IN N DRV 250
C UNKNOWNNS. MATRIX W IS A GIVEN DIAGONAL MATRIX OF WEIGHTS WITH ALL DRV 260

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C DIAGONAL ELEMENTS NONNEGATIVE. LET H = (SQRT(W))*A.          DRV 270
C                                                                DRV 280
C MODE = 1 INDICATES THAT IF N1.LT.N THE ORIGINAL MATRIX H (M BY N)  DRV 290
C IS TO BE REPLACED BY A SMALLER MATRIX (M BY N1) WHOSE          DRV 300
C COLUMNS ARE LINEARLY INDEPENDENT, AND A SOLUTION IS TO BE     DRV 310
C SOUGHT FOR THE SMALLER SYSTEM OF EQUATIONS.  THUS N - N1      DRV 320
C COLUMNS OF THE ORIGINAL MATRIX H ARE DELETED, AND            DRV 330
C COEFFICIENTS CORRESPONDING TO THESE N - N1 DELETED COLUMNS   DRV 340
C WILL BE SET EQUAL TO ZERO.                                     DRV 350
C MODE = 2 INDICATES THAT A SOLUTION IS SOUGHT FOR A LEAST SQUARES  DRV 360
C PROBLEM HAVING N ELEMENTS IN THE SOLUTION VECTOR.  IN ORDER   DRV 370
C TO OBTAIN A UNIQUE SOLUTION, THE CONDITION THAT THE           DRV 380
C SOLUTION VECTOR BE OF MINIMAL EUCLIDEAN NORM IS IMPOSED.      DRV 390
C                                                                DRV 400
C THE SEQUENCE OF INPUT CARDS FOR THIS MAIN PROGRAM IS --       DRV 410
C 0. CARD SPECIFYING MODE DESIRED FOR PROBLEMS WHICH FOLLOW, IN (I5)  DRV 420
C FORMAT.                                                         DRV 430
C 1. PROBLEM HEADING CARD, IN (80A1) FORMAT.                     DRV 440
C 2. PARAMETER CARD IN (6I5,5X,F10.0) FORMAT, GIVING VALUES OF THE  DRV 450
C PARAMETERS M, N, M1, L, ITYPE, IWGHT, TOL.                     DRV 460
C M TOTAL NUMBER OF EQUATIONS.                                   DRV 470
C N NUMBER OF UNKNOWN COEFFICIENTS.                             DRV 480
C M1 NUMBER OF LINEAR CONSTRAINTS (0.LE.M1.LE.M AND M1.LE.N).   DRV 490
C L NUMBER OF RIGHT-HAND SIDES (VECTORS OF OBSERVATIONS).       DRV 500
C ITYPE PARAMETER WHICH SPECIFIES WHETHER OR NOT DATA FOR A     DRV 510
C POLYNOMIAL TYPE FIT ARE TO BE READ IN.                        DRV 520
C ITYPE = 1 INDICATES POLYNOMIAL TYPE.                           DRV 530
C ITYPE = 2 INDICATES NON-POLYNOMIAL TYPE.                       DRV 540
C IWGHT PARAMETER WHICH SPECIFIES WHETHER OR NOT WEIGHTS ARE TO   DRV 550
C BE READ IN.                                                    DRV 560
C IWGHT = 1 INDICATES WEIGHTS ARE NOT TO BE READ IN. (THE        DRV 570
C PROGRAM SETS ALL WEIGHTS EQUAL TO 1.0.)                        DRV 580
C IWGHT = 2 INDICATES WEIGHTS ARE TO BE READ IN.                 DRV 590
C TOL PARAMETER USED IN DETERMINING THE RANK OF MATRIX H.        DRV 600
C NOTE --                                                         DRV 610
C (1) IF TOL EQUALS ZERO, THE TOLERANCE USED IN THE              DRV 620
C DECOMPOSITION SUBROUTINE WILL BE BASED ON MACHINE              DRV 630
C PRECISION.                                                      DRV 640
C (2) IF TOL IS GREATER THAN ZERO, THIS VALUE OF TOL WILL BE   DRV 650
C USED IN SETTING AN ABSOLUTE TOLERANCE FOR COMPARISON          DRV 660
C WITH DIAGONAL ELEMENTS OF THE TRIANGULAR MATRIX OBTAINED     DRV 670
C IN THE DECOMPOSITION SUBROUTINE.  THE VALUE OF TOL CAN        DRV 680
C BE BASED ON KNOWLEDGE CONCERNING THE ACCURACY OF THE          DRV 690
C DATA.                                                          DRV 700
C 3. CARD GIVING FORMAT OF THE DATA CARDS (CONTAINING A, B AND   DRV 710
C POSSIBLY W) WHICH FOLLOW.  THIS FORMAT CARD IS IN (80A1)      DRV 720
C FORMAT.                                                         DRV 730
C 4. DATA CARDS FOR THE ARRAYS A, B AND POSSIBLY W.  THERE ARE FOUR  DRV 740
C POSSIBLE CONFIGURATIONS FOR THE DATA, DEPENDING ON THE VALUES  DRV 750
C OF ITYPE AND IWGHT.  (FOR POLYNOMIAL FITS, THE FIRST POWER OF A  DRV 760
C IS READ IN AND HIGHER POWERS ARE COMPUTED BY THE PROGRAM WHEN  DRV 770
C N.GT.2.)  THE FOUR CONFIGURATIONS ARE ILLUSTRATED BELOW BY    DRV 780
C SHOWING WHAT THE CARD (OR CARDS) FOR THE I-TH ROW OF DATA    DRV 790
C CONTAINS.                                                       DRV 800
C A. ITYPE = 1, IWGHT = 1.                                       DRV 810
C POLYNOMIAL TYPE FIT.  EQUAL WEIGHTS, NOT TO BE READ IN.       DRV 820
C A(I,2) B(I,1) B(I,2) ... B(I,L)                                DRV 830
C B. ITYPE = 1, IWGHT = 2.                                       DRV 840
C POLYNOMIAL TYPE FIT.  UNEQUAL WEIGHTS, TO BE READ IN.        DRV 850
C A(I,2) B(I,1) B(I,2) ... B(I,L) W(I)                          DRV 860
C C. ITYPE = 2, IWGHT = 1.                                       DRV 870
C NON-POLYNOMIAL TYPE FIT.  EQUAL WEIGHTS, NOT TO BE READ IN.   DRV 880
C A(I,1) A(I,2) ... A(I,N) B(I,1) B(I,2) ... B(I,L)             DRV 890
C D. ITYPE = 2, IWGHT = 2.                                       DRV 900
C NON-POLYNOMIAL TYPE FIT.  UNEQUAL WEIGHTS, TO BE READ IN.    DRV 910
C A(I,1) A(I,2) ... A(I,N) B(I,1) B(I,2) ... B(I,L) W(I)       DRV 920
C 5. CARD IN (I5) FORMAT GIVING VALUE OF THE PARAMETER IFDONE.  IF  DRV 930
C THE PROBLEM AT HAND IS TO BE FOLLOWED BY ANOTHER PROBLEM,    DRV 940
C IFDONE = 1.  OTHERWISE, IFDONE EQUALS ANY INTEGER EXCEPT 1.  DRV 950
C                                                                DRV 960
C DIMENSIONS OF ARRAYS ARE SET ASSUMING THAT M.LE.21, N.LE.8, AND  DRV 970
C L.LE.3.                                                         DRV 980
C                                                                DRV 990
C INTEGER IPIVOT(8)                                              DRV 1000
C REAL A(21,8), B(21,3), RES(21,3), TOL, W(21), X(8,3)         DRV 1010
C ARRAYS Q AND R ARE USED IN SUBROUTINE L2A.                     DRV 1020

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C ARRAY QR IS USED IN SUBROUTINE L2B.                                DRV 1030
  REAL Q(21,8), R(8,8)                                             DRV 1040
  REAL QR(29,8)                                                    DRV 1050
C SUBROUTINE L2A REQUIRES THAT C BE DIMENSIONED AT LEAST 4*(M+N) + 2*L. DRV 1060
C SUBROUTINE L2B REQUIRES THAT C BE DIMENSIONED AT LEAST 6*(M+N) + 2*L. DRV 1070
  REAL C(180)                                                       DRV 1080
  REAL SD, SDX(8), SNORM, SS, Z                                     DRV 1090
  LOGICAL FAIL(3)                                                  DRV 1100
  DIMENSION HEAD(80), IFMT(80)                                     DRV 1110
C                                                                      DRV 1120
C IN THE FOLLOWING DATA STATEMENT, NR IS THE CARD READER DEVICE   DRV 1130
C AND NW IS THE PRINTER DEVICE NUMBER.                             DRV 1140
C                                                                      DRV 1150
  DATA NR,NW /5,6/                                               DRV 1160
C                                                                      DRV 1170
C THE PARAMETERS MM, NN AND MMPNN WHICH APPEAR IN THE STATEMENTS  DRV 1180
C CALL L2A(...) AND CALL L2B(...) ARE USED IN SETTING ADJUSTABLE DRV 1190
C DIMENSIONS OF ARRAYS. THEY ARE GIVEN SPECIFIC VALUES IN THE   DRV 1200
C FOLLOWING DATA STATEMENTS. MMPNN IS USED ONLY IN CALL L2B(...). DRV 1210
C                                                                      DRV 1220
  DATA MM,NN /21,8/                                              DRV 1230
  DATA MMPNN /29/                                               DRV 1240
C                                                                      DRV 1250
  READ (NR,99730) MODE                                           DRV 1260
C                                                                      DRV 1270
C DEFAULT VALUE FOR MODE IS 1.                                     DRV 1280
C                                                                      DRV 1290
  IF (MODE.NE.2) MODE = 1                                         DRV 1300
  IPROB = 0                                                         DRV 1310
10 READ (NR,99990) HEAD                                           DRV 1320
  IPROB = IPROB + 1                                               DRV 1330
  WRITE (NW,99980) IPROB                                          DRV 1340
  WRITE (NW,99970) HEAD                                           DRV 1350
  READ (NR,99960) M,N,M1,L,ITYPE,IWGHT,TOL                       DRV 1360
  WRITE (NW,99950)                                               DRV 1370
  WRITE (NW,99940) M,N,M1,L,ITYPE,IWGHT,MODE,TOL                DRV 1380
  READ (NR,99990) IFMT                                           DRV 1390
  WRITE (NW,99930) IFMT                                           DRV 1400
  GO TO (20,100), ITYPE                                           DRV 1410
C                                                                      DRV 1420
C TYPE 1. POLYNOMIAL FIT.                                         DRV 1430
20 GO TO (30,50), IWGHT                                           DRV 1440
C   A. EQUAL WEIGHTS, NOT TO BE READ IN.                          DRV 1450
30 DO 40 I=1,M                                                    DRV 1460
  READ (NR,IFMT) A(I,2), (B(I,K),K=1,L)                          DRV 1470
  W(I) = 1.0                                                       DRV 1480
40 CONTINUE                                                       DRV 1490
  GO TO 70                                                         DRV 1500
C   B. UNEQUAL WEIGHTS, TO BE READ IN.                            DRV 1510
50 DO 60 I=1,M                                                    DRV 1520
  READ (NR,IFMT) A(I,2), (B(I,K),K=1,L),W(I)                    DRV 1530
60 CONTINUE                                                       DRV 1540
70 DO 90 I=1,M                                                    DRV 1550
  A(I,1) = 1.0                                                    DRV 1560
  IF (N.LT.3) GO TO 90                                           DRV 1570
  DO 80 J=3,N                                                      DRV 1580
  A(I,J) = A(I,2)**(J-1)                                         DRV 1590
80 CONTINUE                                                       DRV 1600
90 CONTINUE                                                       DRV 1610
  GO TO 150                                                        DRV 1620
C                                                                      DRV 1630
C TYPE 2. NON-POLYNOMIAL FIT.                                     DRV 1640
100 GO TO (110,130), IWGHT                                        DRV 1650
C   C. EQUAL WEIGHTS, NOT TO BE READ IN.                          DRV 1660
110 DO 120 I=1,M                                                  DRV 1670
  READ (NR,IFMT) (A(I,J),J=1,N), (B(I,K),K=1,L)                 DRV 1680
  W(I) = 1.0                                                       DRV 1690
120 CONTINUE                                                       DRV 1700
  GO TO 150                                                        DRV 1710
C   D. UNEQUAL WEIGHTS, TO BE READ IN.                            DRV 1720
130 DO 140 I=1,M                                                  DRV 1730
  READ (NR,IFMT) (A(I,J),J=1,N), (B(I,K),K=1,L),W(I)           DRV 1740
140 CONTINUE                                                       DRV 1750
150 IF (M1.EQ.0 .OR. IWGHT.EQ.1) GO TO 170                       DRV 1760
C                                                                      DRV 1770
C INSURE THAT WEIGHTS EQUAL 1.0 FOR THE FIRST M1 EQUATIONS WHEN M1 DRV 1780

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C IS GREATER THAN ZERO.                                DRV 1790
C                                                        DRV 1800
    DO 160 I=1,M1                                       DRV 1810
      W(I) = 1.0                                         DRV 1820
    160 CONTINUE                                       DRV 1830
C                                                        DRV 1840
C PRINT A, B AND W.                                    DRV 1850
C                                                        DRV 1860
    170 WRITE (NW,99920)                                 DRV 1870
      KZ = 0                                             DRV 1880
      DO 180 I=1,M                                       DRV 1890
        IF (W(I).EQ.0.0) KZ = KZ + 1                   DRV 1900
        WRITE (NW,99910) (A(I,J),J=1,N),(B(I,K),K=1,L),W(I) DRV 1910
    180 CONTINUE                                       DRV 1920
C                                                        DRV 1930
    GO TO (190,200), MODE                               DRV 1940
C                                                        DRV 1950
    190 CALL L2A(M, N, M1, L, A, B, W, TOL, MM, NN,     DRV 1960
      * N1, IPIVOT, X, RES, R, Q, C, IFAULT)           DRV 1970
C                                                        DRV 1980
    GO TO 210                                           DRV 1990
C                                                        DRV 2000
    200 CALL L2B(M, N, M1, L, A, B, W, TOL, MM, NN, MMPNN, DRV 2010
      * N1, IPIVOT, X, RES, QR, C, IFAULT)             DRV 2020
C                                                        DRV 2030
C PRINT COMPUTED RESULTS.                               DRV 2040
C                                                        DRV 2050
    210 WRITE (NW,99900)                                 DRV 2060
      WRITE (NW,99890) MODE,IFAU
      IF (IFAU
      WRITE (NW,99880) N1
      IF (N1.EQ.0) GO TO 390
      IF (N1.LT.M1) GO TO 390
      WRITE (NW,99870)
      WRITE (NW,99860) (IPIVOT(J),J=1,N1)
      IF (N1.EQ.N) GO TO 220
      N1P1 = N1 + 1
      WRITE (NW,99850)
      WRITE (NW,99860) (IPIVOT(J),J=N1P1,N)
    220 NDF = M - N1 - KZ
      WRITE (NW,99720) KZ,NDF
      WRITE (NW,99840)
      DO 240 K=1,L
        K2 = L + K
        IF (C(K).LT.0.0) GO TO 230
        FAIL(K) = .FALSE.
        WRITE (NW,99830) K,C(K),C(K2)
        GO TO 240
    230 FAIL(K) = .TRUE.
        C(K) = -C(K)
        WRITE (NW,99820) K,C(K),C(K2)
    240 CONTINUE
      IF (IFAU
      DO 350 K=1,L
C
C COMPUTE SUM OF SQUARED RESIDUALS, NORM OF RESIDUALS, RESIDUAL
C STANDARD DEVIATION, STANDARD DEVIATIONS OF COEFFICIENTS, AND
C PREDICTED VALUES. PRINT THESE QUANTITIES, TOGETHER WITH
C COEFFICIENTS, OBSERVED VALUES AND RESIDUALS.
C
      IF (FAIL(K)) GO TO 350
      WRITE (NW,99810) K
      SS = 0.0
      IF (M.EQ.M1) GO TO 310
      M1P1 = M1 + 1
      DO 250 I=M1P1,M
        SS = SS + (RES(I,K)*W(I))**2
    250 CONTINUE
      IF (M.LE.N1+KZ) GO TO 310
      IF (MODE.EQ.2 .AND. N1.LT.N) GO TO 310
      SD = SQRT(SS/FLOAT(NDF))
      IF (MODE.EQ.2) GO TO 270
      DO 260 J=1,N
        IF (R(J,J).LT.0.0) R(J,J) = 0.0
        SDX(J) = SD*SQRT(R(J,J))
    260 CONTINUE

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      GO TO 290
270 DO 280 J=1,N
      IF (QR(J,J).LT.0.0) QR(J,J) = 0.0
      SDX(J) = SD*SQRT(QR(J,J))
280 CONTINUE
290 WRITE (NW,99800)
      DO 300 J=1,N
        WRITE (NW,99770) J,X(J,K),SDX(J)
300 CONTINUE
      GO TO 330
310 WRITE (NW,99790)
      DO 320 J=1,N
        WRITE (NW,99770) J,X(J,K)
320 CONTINUE
330 WRITE (NW,99780)
      DO 340 I=1,M
        Z = B(I,K) - RES(I,K)
        WRITE (NW,99770) I,B(I,K),Z,RES(I,K)
340 CONTINUE
      SNORM = SQRT(SS)
      WRITE (NW,99760) SS,SNORM
      IF ((MODE.EQ.2 .AND. N1.LT.N) .OR. (M.EQ.N1+KZ)) GO TO 350
      WRITE (NW,99750) SD
350 CONTINUE
C
C PRINT LOWER TRIANGULAR PORTION OF SYMMETRIC UNSCALED COVARIANCE
C MATRIX.
C
      IF (MODE.EQ.2 .AND. N1.LT.N) GO TO 390
      WRITE (NW,99740)
      IF (MODE.EQ.2) GO TO 370
      DO 360 I=1,N
        WRITE (NW,99910) (R(I,J),J=1,I)
360 CONTINUE
      GO TO 390
370 DO 380 I=1,N
        WRITE (NW,99910) (QR(I,J),J=1,I)
380 CONTINUE
390 READ (NR,99730) IFDONE
      IF (IFDONE.EQ.1) GO TO 10
C
      STOP
C
C FORMAT STATEMENTS.
C
99990 FORMAT (80A1)
99980 FORMAT (1H1,115(1H*),4X,7HPROBLEM,I4)
99970 FORMAT (1H0,80A1)
99960 FORMAT (6I5,5X,F10.0)
99950 FORMAT (1H0,3X,1HM,4X,1HN,3X,2HML,4X,1HL,5X,5HITYPE,5X,4HIWGH,
* 1HT,6X,4HMODE,7X,3HTOL)
99940 FORMAT (4I5,3I10,G15.8)
99930 FORMAT (8H0FORMAT ,80A1)
99920 FORMAT (41H0MATRIX A, MATRIX B AND VECTOR OF WEIGHTS/)
99910 FORMAT (1X,8G15.8)
99900 FORMAT (17H0COMPUTED RESULTS)
99890 FORMAT (7H0MODE =,I4,5X,8HIFAUULT =,I4)
99880 FORMAT (44H0N1 = COMPUTED RANK OF SYSTEM OF EQUATIONS =,I4)
99870 FORMAT (50H0COLUMNS OF H = (SQRT(W))*A WERE SELECTED BY THE P,
* 37HIVOTING SCHEME IN THE FOLLOWING ORDER/)
99860 FORMAT (30I4)
99850 FORMAT (50H0THE FOLLOWING COLUMNS OF H ARE LINEARLY DEPENDENT,
* 48H. IF MODE 1, THEY DID NOT ENTER THE REGRESSION./6H IF MO,
* 24HDE 2, THEY ENTERED LAST./)
99840 FORMAT (1H0,15X,9HREPORT ON,12X,9HNUMBER OF,4X,11HESTIMATED N,
* 16HNUMBER OF CORRECT/13H B-VECTOR NO.,3X,11HCONVERGENCE,10X,
* 10HITERATIONS,3X,26HDIGITS IN INITIAL SOLUTION/)
99830 FORMAT (17,9X,9HCONVERGED,14X,F4.0,10X,G15.8)
99820 FORMAT (17,9X,18HFAILED TO CONVERGE,5X,F4.0,10X,G15.8)
99810 FORMAT (26H0SOLUTION FOR B-VECTOR NO.,I3)
99800 FORMAT (1H0,27X,18HSTANDARD DEVIATION/5X,1HJ,4X,10HCOEFFICIENT,
* 4HT(J),5X,17HOF COEFFICIENT(J)/)
99790 FORMAT (1H0,4X,1HJ,4X,14HCOEFFICIENT(J)/)
99780 FORMAT (1H0,4X,1HI,4X,11HOBSERVED(I),9X,12HPREDICTED(I),10X,
* 11HRESIDUAL(I)/)
99770 FORMAT (I6,G17.8,G21.8)

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DRV 2550
DRV 2560
DRV 2570
DRV 2580
DRV 2590
DRV 2600
DRV 2610
DRV 2620
DRV 2630
DRV 2640
DRV 2650
DRV 2660
DRV 2670
DRV 2680
DRV 2690
DRV 2700
DRV 2710
DRV 2720
DRV 2730
DRV 2740
DRV 2750
DRV 2760
DRV 2770
DRV 2780
DRV 2790
DRV 2800
DRV 2810
DRV 2820
DRV 2830
DRV 2840
DRV 2850
DRV 2860
DRV 2870
DRV 2880
DRV 2890
DRV 2900
DRV 2910
DRV 2920
DRV 2930
DRV 2940
DRV 2950
DRV 2960
DRV 2970
DRV 2980
DRV 2990
DRV 3000
DRV 3010
DRV 3020
DRV 3030
DRV 3040
DRV 3050
DRV 3060
DRV 3070
DRV 3080
DRV 3090
DRV 3100
DRV 3110
DRV 3120
DRV 3130
DRV 3140
DRV 3150
DRV 3160
DRV 3170
DRV 3180
DRV 3190
DRV 3200
DRV 3210
DRV 3220
DRV 3230
DRV 3240
DRV 3250
DRV 3260
DRV 3270
DRV 3280
DRV 3290
DRV 3300

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99760 FORMAT (30H0SUM OF SQUARED RESIDUALS =,G15.8/1X,8HNORM OF , DRV 3310
* 9HRESIDUALS,11X,1H=,G15.8) DRV 3320
99750 FORMAT (30H RESIDUAL STANDARD DEVIATION =,G15.8) DRV 3330
99740 FORMAT (27H0UNSCALED COVARIANCE MATRIX/) DRV 3340
99730 FORMAT (15) DRV 3350
99720 FORMAT(25H0NUMBER OF ZERO WEIGHTS =,I3,5X,17HDEG. OF FREEDOM =,I3) DRV 3360
END DRV 3370

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SUBROUTINE L2A(M, N, M1, L, A, B, W, TOL, MM, NN, L2A 10
* N1, IPIVOT, X, RES, R, Q, C, IFAULT) L2A 20
C ** PURPOSE ** L2A 30
C SUBROUTINE L2A COMPUTES LEAST SQUARES SOLUTIONS TO OVERDETERMINED L2A 40
C SYSTEMS OF LINEAR EQUATIONS. THE METHOD USED IS A MODIFIED L2A 50
C GRAM-SCHMIDT ORTHOGONAL DECOMPOSITION WITH ITERATIVE REFINEMENT OF L2A 60
C THE SOLUTION. THE SOLUTION MAY BE SUBJECT TO LINEAR EQUALITY L2A 70
C CONSTRAINTS. OUTPUT INCLUDES THE LEAST SQUARES COEFFICIENTS, L2A 80
C RESIDUALS, UNSCALED COVARIANCE MATRIX, AND INFORMATION ON THE L2A 90
C BEHAVIOR OF THE ITERATIVE REFINEMENT PROCEDURE. L2A 100
C MATRIX A IS THE GIVEN MATRIX OF A SYSTEM OF M LINEAR EQUATIONS IN N L2A 110
C UNKNOWN, AND MATRIX W IS A GIVEN DIAGONAL MATRIX OF WEIGHTS WITH ALL L2A 120
C DIAGONAL ELEMENTS NONNEGATIVE. LET H = (SQRT(W))*A. L2A 130
C IN THE EVENT THAT N1 (THE COMPUTED RANK OF MATRIX H) IS LESS THAN N L2A 140
C (THE NUMBER OF UNKNOWN COEFFICIENTS), THE ORIGINAL MATRIX H (M BY N) L2A 150
C IS REPLACED BY A SMALLER MATRIX (M BY N1) WHOSE COLUMNS ARE LINEARLY L2A 160
C INDEPENDENT, AND A SOLUTION IS SOUGHT FOR THE SMALLER SYSTEM OF L2A 170
C EQUATIONS. THUS N - N1 COLUMNS OF THE ORIGINAL MATRIX H ARE DELETED, L2A 180
C AND COEFFICIENTS CORRESPONDING TO THESE N - N1 DELETED COLUMNS WILL L2A 190
C BE SET EQUAL TO ZERO. L2A 200
C L2A 210
C ** INPUT VARIABLES ** L2A 220
C M TOTAL NUMBER OF EQUATIONS. L2A 230
C N NUMBER OF UNKNOWN COEFFICIENTS. L2A 240
C M1 NUMBER OF LINEAR CONSTRAINTS (0.LE.M1.LE.M AND M1.LE.N). L2A 250
C L NUMBER OF RIGHT-HAND SIDES (VECTORS OF OBSERVATIONS). L2A 260
C A TWO-DIMENSIONAL ARRAY OF SIZE (MM,N). ON ENTRY, THE ARRAY A L2A 270
C CONTAINS THE GIVEN MATRIX OF A SYSTEM OF M LINEAR EQUATIONS L2A 280
C IN N UNKNOWN, WHERE THE FIRST M1 EQUATIONS ARE TO BE L2A 290
C SATISFIED EXACTLY. A IS LEFT INTACT ON EXIT. L2A 300
C B TWO-DIMENSIONAL ARRAY OF SIZE (MM,L). ON ENTRY, B CONTAINS L2A 310
C THE L GIVEN RIGHT-HAND SIDES (VECTORS OF OBSERVATIONS). B IS L2A 320
C LEFT INTACT ON EXIT. L2A 330
C W VECTOR OF SIZE M. ON ENTRY, W CONTAINS THE DIAGONAL ELEMENTS L2A 340
C OF A GIVEN DIAGONAL MATRIX OF WEIGHTS, ALL NONNEGATIVE. L2A 350
C (THE FIRST M1 ELEMENTS OF W ARE SET EQUAL TO 1.0 BY THE L2A 360
C PROGRAM WHEN M1 IS GREATER THAN ZERO.) ON EXIT, THE ORIGINAL L2A 370
C ELEMENTS OF W HAVE BEEN REPLACED BY THEIR SQUARE ROOTS. L2A 380
C TOL PARAMETER USED IN DETERMINING THE RANK OF MATRIX H. L2A 390
C NOTE -- L2A 400
C (1) IF TOL EQUALS ZERO, THE TOLERANCE USED IN SUBROUTINE L2A 410
C DECOM1 WILL BE BASED ON MACHINE PRECISION. L2A 420
C (2) IF TOL IS GREATER THAN ZERO, THIS VALUE OF TOL WILL BE L2A 430
C USED IN SETTING AN ABSOLUTE TOLERANCE FOR COMPARISON WITH L2A 440
C DIAGONAL ELEMENTS OF THE TRIANGULAR MATRIX OBTAINED IN L2A 450
C SUBROUTINE DECOM1. THE VALUE OF TOL CAN BE BASED ON L2A 460
C KNOWLEDGE CONCERNING THE ACCURACY OF THE DATA. L2A 470
C MM DIMENSIONING PARAMETER SPECIFYING MAXIMUM NUMBER OF ROWS IN L2A 480
C THE ARRAYS A, B, RES AND Q. MM MUST SATISFY MM.GE.M. L2A 490
C NN DIMENSIONING PARAMETER SPECIFYING MAXIMUM NUMBER OF ROWS IN L2A 500
C THE ARRAYS X AND R. NN MUST SATISFY NN.GE.N. L2A 510
C L2A 520
C ** OUTPUT VARIABLES AND INTERNAL VARIABLES ** L2A 530
C N1 COMPUTED RANK OF MATRIX H, WHERE H = (SQRT(W))*A. L2A 540
C IPIVOT VECTOR OF SIZE N. ON EXIT, THIS ARRAY RECORDS THE ORDER L2A 550
C IN WHICH THE COLUMNS OF H WERE SELECTED BY THE PIVOTING L2A 560
C SCHEME IN THE COURSE OF THE ORTHOGONAL DECOMPOSITION. L2A 570
C WHENEVER N1.LT.N, THE FIRST N1 ELEMENTS OF IPIVOT INDICATE L2A 580
C WHICH COLUMNS OF H WERE FOUND TO BE LINEARLY INDEPENDENT. L2A 590
C X TWO-DIMENSIONAL ARRAY OF SIZE (NN,L). ON EXIT, X CONTAINS L2A 600
C THE SOLUTION VECTORS. L2A 610
C RES TWO-DIMENSIONAL ARRAY OF SIZE (MM,L). ON EXIT, RES CONTAINS L2A 620
C THE RESIDUAL VECTORS. L2A 630
C R TWO-DIMENSIONAL ARRAY OF SIZE (NN,N). ON EXIT, R CONTAINS L2A 640
C THE LOWER TRIANGULAR PORTION OF THE SYMMETRIC UNSCALED L2A 650
C COVARIANCE MATRIX. (THIS ARRAY IS USED INTERNALLY TO STORE L2A 660

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C      RESULTS FROM SUBROUTINE DECOM1 WHICH ARE DESTROYED IN          L2A 670
C      COMPUTING THE COVARIANCE MATRIX.)                             L2A 680
C Q    TWO-DIMENSIONAL ARRAY OF SIZE (MM,N) USED INTERNALLY ONLY.   L2A 690
C C    VECTOR HAVING AT LEAST 4*(M+N)+2*L ELEMENTS USED (1) FOR      L2A 700
C      INTERNAL WORK SPACE AND (2) FOR RETURNING INFORMATION ON THE   L2A 710
C      BEHAVIOR OF THE ITERATIVE REFINEMENT PROCEDURE.              L2A 720
C      (A) NUMIT IS THE NUMBER OF ITERATIONS CARRIED OUT DURING THE  L2A 730
C          ITERATIVE REFINEMENT IN ATTEMPTING TO OBTAIN A SOLUTION   L2A 740
C          FOR THE K-TH RIGHT-HAND SIDE.                             L2A 750
C          ON EXIT, C(K) = +NUMIT IF THE SOLUTION CONVERGED, AND      L2A 760
C          C(K) = -NUMIT IF THE SOLUTION FAILED TO CONVERGE.        L2A 770
C      (B) DIGITX GIVES AN ESTIMATE OF THE NUMBER OF CORRECT DIGITS  L2A 780
C          IN THE INITIAL SOLUTION OF THE COEFFICIENTS FOR THE K-TH L2A 790
C          RIGHT-HAND SIDE. ON EXIT, C(K+L) = DIGITX.                L2A 800
C IFAULT FAULT INDICATOR WHICH IS ZERO IF NO ERRORS WERE ENCOUNTERED L2A 810
C      AND POSITIVE IF ERRORS WERE DETECTED OR IF EVIDENCE OF SEVERE L2A 820
C      ILL-CONDITIONING WAS FOUND. DIAGNOSTIC MESSAGES ARE PRINTED  L2A 830
C      FROM SUBROUTINE ERROR. IF IFAULT IS SET EQUAL TO 1, 2, 3, 4,  L2A 840
C      5, 6 OR 7, EXECUTION IS TERMINATED. EXECUTION CONTINUES WHEN L2A 850
C      IFAULT IS SET EQUAL TO 8, 9 OR 10 PROVIDED THAT A SOLUTION    L2A 860
C      WAS OBTAINED FOR AT LEAST ONE RIGHT-HAND SIDE. THE VALUE OF  L2A 870
C      IFAULT IS USED TO INDICATE THE FOLLOWING --                  L2A 880
C      0 = NO ERRORS ENCOUNTERED.                                    L2A 890
C      1 = BAD INPUT PARAMETER (M, N OR L).                          L2A 900
C      2 = BAD INPUT PARAMETER (M1).                                 L2A 910
C      3 = BAD DIMENSION. EITHER M.GT.MM OR N.GT.NN.                L2A 920
C      4 = AT LEAST ONE WEIGHT IS NEGATIVE.                          L2A 930
C      5 = EITHER MATRIX H OR MATRIX OF CONSTRAINTS EQUALS ZERO.    L2A 940
C      6 = CONSTRAINTS ARE LINEARLY DEPENDENT.                       L2A 950
C      7 = ALL SOLUTIONS FAILED TO CONVERGE.                         L2A 960
C      8 = SOLUTION FAILED TO CONVERGE FOR AT LEAST ONE RIGHT-HAND  L2A 970
C          SIDE.                                                    L2A 980
C      9 = LARGE NUMBER OF ITERATIONS REQUIRED FOR CONVERGENCE.      L2A 990
C      10 = ESTIMATED NUMBER OF DIGITS IN INITIAL SOLUTION OF        L2A 1000
C          COEFFICIENTS IS SMALL.                                    L2A 1010
C      11 = DIAGONAL ELEMENT OF COVARIANCE MATRIX WAS COMPUTED TO BE L2A 1020
C          NEGATIVE OWING TO ROUNDING ERROR.                         L2A 1030
C
C ** SUBROUTINES REQUIRED **                                         L2A 1040
C SUBROUTINE DECOM1                                                 L2A 1050
C      USES MODIFIED GRAM-SCHMIDT ALGORITHM WITH PIVOTING TO        L2A 1060
C      OBTAIN AN ORTHOGONAL DECOMPOSITION OF THE INPUT MATRIX.      L2A 1070
C SUBROUTINE SOLVE1                                                 L2A 1080
C      COMPUTES COEFFICIENTS AND RESIDUALS. ITERATIVE REFINEMENT IS L2A 1090
C      USED TO IMPROVE THE ACCURACY OF THE INITIAL SOLUTION.        L2A 1100
C SUBROUTINE COVAR                                                 L2A 1110
C      COMPUTES UNSCALED COVARIANCE MATRIX OF THE COEFFICIENTS.    L2A 1120
C SUBROUTINE ERROR                                                 L2A 1130
C      PRINTS ERROR DIAGNOSTICS WHEN ERRORS ARE DETECTED OR WHEN   L2A 1140
C      EVIDENCE OF SEVERE ILL-CONDITIONING IS FOUND.               L2A 1150
C
C ** STORAGE REQUIREMENTS **                                       L2A 1160
C THE STORAGE REQUIRED FOR THE DIMENSIONED ARRAYS IN SUBROUTINE L2A IS L2A 1170
C      M*(2*N + 2*L + 5) + N*(N + L + 5) + 2*L                      L2A 1180
C LOCATIONS. ALL ARRAYS REQUIRED IN SUBROUTINES CALLED BY L2A ARE   L2A 1190
C      DECLARED HEREIN AND ARE TRANSMITTED ONLY THROUGH PARAMETER LISTS OF L2A 1200
C      CALL-SEQUENCES.                                             L2A 1210
C
C ** PRECISION OF ARITHMETIC CALCULATIONS **                       L2A 1220
C SINGLE PRECISION ARITHMETIC IS USED FOR ALL CALCULATIONS EXCEPT L2A 1230
C DOUBLE PRECISION ACCUMULATION OF INNER PRODUCTS. (THE VARIABLE SUM L2A 1240
C IS DECLARED TO BE DOUBLE PRECISION IN SUBROUTINES DECOM1, SOLVE1 AND L2A 1250
C COVAR.) IT IS ESSENTIAL FOR THE SUCCESS OF THE ITERATIVE REFINEMENT L2A 1260
C PROCEDURE THAT INNER PRODUCTS BE ACCUMULATED IN DOUBLE PRECISION. L2A 1270
C
C ** CONVERSION OF THE PROGRAM TO DOUBLE PRECISION **              L2A 1280
C *****                                                           L2A 1290
C * ON COMPUTERS HAVING SHORT WORD LENGTH (AS THE IBM 360/370) IT MAY * L2A 1300
C * BE DESIRABLE TO PERFORM ALL CALCULATIONS IN DOUBLE PRECISION. IN * L2A 1310
C * THIS CASE, THE ITERATIVE REFINEMENT PRESENTLY INCLUDED IN SOLVE1 * L2A 1320
C * SHOULD BE OMITTED.                                             * L2A 1330
C * TO CONVERT THE PROGRAM TO DOUBLE PRECISION, THE FOLLOWING      * L2A 1340
C * APPROACH IS SUGGESTED.                                         * L2A 1350
C *
C * 1. VARIABLES PRESENTLY DECLARED TO BE REAL SHOULD BE DECLARED * L2A 1360

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C * DOUBLE PRECISION. THOSE TYPED INTEGER, DOUBLE PRECISION AND * L2A 1420
C * LOGICAL SHOULD NOT BE CHANGED. * L2A 1430
C * 2. THE USE OF FAIL, NUMIT AND DIGITX SHOULD BE OMITTED. * L2A 1440
C * 3. DESCRIPTION OF VARIABLE C (AT L2A 700-800) SHOULD READ -- * L2A 1450
C * C VECTOR HAVING AT LEAST 4*(M+N) ELEMENTS USED ONLY FOR * L2A 1460
C * INTERNAL WORK SPACE. * L2A 1470
C * 4. THE VALUE OF ETA (AT L2A 1930) SHOULD BE SET SO THAT IT IS THE * L2A 1480
C * SMALLEST POSITIVE DOUBLE PRECISION NUMBER SUCH THAT 1.0 + ETA * L2A 1490
C * IS GREATER THAN 1.0 IN DOUBLE PRECISION ARITHMETIC. * L2A 1500
C * FOR IBM COMPUTER TYPE, ETA = 16.**(-13) * L2A 1510
C * FOR UNIVAC COMPUTER TYPE, ETA = 2.**(-59) * L2A 1520
C * 5. THE FOLLOWING FORTRAN FUNCTIONS SHOULD BE CHANGED -- * L2A 1530
C * SINGLE PRECISION NAME DOUBLE PRECISION NAME * L2A 1540
C * DBLE(X) X * L2A 1550
C * FLOAT(N) DBLE(FLOAT(N)) * L2A 1560
C * Sqrt(X) DSqrt(X) * L2A 1570
C * DBLE(X) IS USED IN SUBROUTINES DECOM1, SOLVE1 AND COVAR. * L2A 1580
C * FLOAT(N) IS USED IN SUBROUTINE DECOM1. * L2A 1590
C * Sqrt(X) IS USED IN SUBROUTINE L2A. * L2A 1600
C * 6. IT MAY BE NECESSARY OR DESIRABLE TO CHANGE CERTAIN FORMATS IN * L2A 1610
C * SUBROUTINE ERROR, REPLACING G SPECIFICATIONS BY D. * L2A 1620
C * 7. REPLACE STATEMENT L2A 2450 BY A STATEMENT READING * L2A 1630
C * K3 = 1 * L2A 1640
C * 8. FURTHER DETAILS ARE GIVEN IN SUBROUTINE SOLVE1 IN CONNECTION * L2A 1650
C * WITH THE OMISSION OF ITERATIVE REFINEMENT. * L2A 1660
C * 9. IN SUBROUTINE L2A, STATEMENTS L2A 960-1010, 1790-1800, 1990, * L2A 1670
C * 2320-2330, 2430-2440, 2970, 3170-3460 AND 3480-3510 SHOULD BE * L2A 1680
C * OMITTED. * L2A 1690
C * STATEMENT NUMBERS GIVEN HERE REFER TO THOSE IN THE RIGHT-HAND * L2A 1700
C * MARGIN. CERTAIN COMMENTS IN SUBROUTINE L2A DO NOT APPLY TO * L2A 1710
C * THE DOUBLE PRECISION VERSION. * L2A 1720
C * * L2A 1730
C ***** L2A 1740
C L2A 1750
C INTEGER IPIVOT(N) L2A 1760
C REAL A(MM,N), B(MM,L), C(1), ETA, Q(MM,N), R(NN,N), L2A 1770
C * RES(MM,L), TOL, W(M), X(NN,L), Z L2A 1780
C REAL DIGITX L2A 1790
C LOGICAL FAIL L2A 1800
C LOGICAL SING L2A 1810
C L2A 1820
C SET VALUE OF ETA, A MACHINE-DEPENDENT PARAMETER. L2A 1830
C ETA, THE RELATIVE MACHINE PRECISION, IS THE SMALLEST POSITIVE REAL L2A 1840
C NUMBER SUCH THAT 1.0 + ETA IS GREATER THAN 1.0 IN FLOATING-POINT L2A 1850
C ARITHMETIC. L2A 1860
C L2A 1870
C FOR IBM COMPUTER TYPE, ETA = 16.**(-5) L2A 1880
C FOR UNIVAC COMPUTER TYPE, ETA = 2.**(-26) L2A 1890
C FOR CDC COMPUTER TYPE, ETA = 2.**(-47) L2A 1900
C FOR HONEYWELL COMPUTER TYPE, ETA = 2.**(-27) L2A 1910
C L2A 1920
C ETA = 2.**(-26) L2A 1930
C L2A 1940
C DEFAULT VALUE FOR TOL IS ZERO. L2A 1950
C L2A 1960
C IF (TOL.LT.0.0) TOL = 0.0 L2A 1970
C IFAULT = 0 L2A 1980
C KSUM = 0 L2A 1990
C L2A 2000
C PERFORM INITIAL CHECKING OF INPUT PARAMETERS, DIMENSIONS AND L2A 2010
C WEIGHTS FOR POSSIBLE ERRORS. L2A 2020
C L2A 2030
C IF (M.GT.0 .AND. N.GT.0 .AND. L.GT.0) GO TO 10 L2A 2040
C IFAULT = 1 L2A 2050
C CALL ERROR(IFAULT, K, Z) L2A 2060
C RETURN L2A 2070
C 10 IF (M.LE.M .AND. M1.LE.N .AND. M1.GE.0) GO TO 20 L2A 2080
C IFAULT = 2 L2A 2090
C CALL ERROR(IFAULT, K, Z) L2A 2100
C RETURN L2A 2110
C 20 IF (M.LE.MM .AND. N.LE.NN) GO TO 30 L2A 2120
C IFAULT = 3 L2A 2130
C K = 1 L2A 2140
C CALL ERROR(IFAULT, K, Z) L2A 2150
C RETURN L2A 2160
C 30 DO 40 I=1,M L2A 2170

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120 CONTINUE
C
CALL SOLVE1(M, N, M1, A, C(K4), W, N1, IPIVOT, Q, R, C(K3),
* ETA, FAIL, NUMIT, DIGITX,
* C(K5), C(K6), C(K7), C(K8), C(K9), C(K10), MM, NN)
C
K0 = K5 - 1
DO 130 J=1,N
K0 = K0 + 1
X(J,K) = C(K0)
130 CONTINUE
IF (M1.EQ.0) GO TO 150
DO 140 I=1,M1
RES(I,K) = 0.0
140 CONTINUE
150 M1P1 = M1 + 1
IF (M1P1.GT.M) GO TO 170
K0 = K6 + M1 - 1
DO 160 I=M1P1,M
K0 = K0 + 1
RES(I,K) = C(K0)
160 CONTINUE
170 CONTINUE
C
C FOR RIGHT-HAND SIDES WHERE CONVERGENCE OF A SOLUTION IS REPORTED,
C A CHECK IS MADE FOR EVIDENCE OF SEVERE ILL-CONDITIONING. SUCH
C EVIDENCE IS FURNISHED BY LARGE VALUES OF NUMIT (NUMBER OF ITERATIONS
C BEFORE CONVERGENCE WAS OBTAINED) AND SMALL VALUES OF DIGITX
C (ESTIMATE OF THE NUMBER OF CORRECT DIGITS IN THE INITIAL SOLUTION
C OF THE COEFFICIENTS). IF NUMIT EXCEEDS -ALOG10(ETA) A DIAGNOSTIC
C MESSAGE IS PRINTED TO WARN OF ILL-CONDITIONING. IF DIGITX IS LESS
C THAN 0.5 (HALF A DECIMAL DIGIT) A SIMILAR DIAGNOSTIC MESSAGE IS
C PRINTED.
C
C(K) = FLOAT(NUMIT)
IF (FAIL) C(K) = -C(K)
K0 = K2 + K - 1
C(K0) = DIGITX
IF (.NOT.FAIL) GO TO 180
C KSUM IS A TALLY OF SOLUTIONS WHICH FAILED TO CONVERGE.
KSUM = KSUM + 1
IFAILT = 8
CALL ERROR(IFAILT, K, Z)
GO TO 200
180 Z = -ALOG10(ETA)
IF (FLOAT(NUMIT).LE.Z) GO TO 190
IFAILT = 9
Z = FLOAT(NUMIT)
CALL ERROR(IFAILT, K, Z)
190 IF (DIGITX.GE.0.5) GO TO 200
IFAILT = 10
Z = DIGITX
CALL ERROR(IFAILT, K, Z)
200 CONTINUE
IF (KSUM.LT.L) GO TO 210
IFAILT = 7
CALL ERROR(IFAILT, K, Z)
RETURN
C
C COMPUTE THE UNSCALED COVARIANCE MATRIX OF THE COEFFICIENTS.
C
210 CALL COVAR(N, M1, N1, IPIVOT, R, C(K3), C(K9), NN)
C
C IN CERTAIN PROBLEMS, SOME DIAGONAL TERMS OF THE UNSCALED COVARIANCE
C MATRIX ARE EQUAL TO ZERO OR TO SMALL POSITIVE NUMBERS. BECAUSE OF
C ROUNDING ERRORS, COMPUTED VALUES FOR THESE TERMS MAY BE SMALL
C NEGATIVE NUMBERS. AN ERROR DIAGNOSTIC IS PRINTED IF ANY DIAGONAL
C TERM IS NEGATIVE.
C
DO 220 J=1,N
IF (R(J,J).GE.0.0) GO TO 220
IFAILT = 11
Z = R(J,J)
CALL ERROR(IFAILT, J, Z)
220 CONTINUE
RETURN
END

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L2A 2940
L2A 2950
L2A 2960
L2A 2970
L2A 2980
L2A 2990
L2A 3000
L2A 3010
L2A 3020
L2A 3030
L2A 3040
L2A 3050
L2A 3060
L2A 3070
L2A 3080
L2A 3090
L2A 3100
L2A 3110
L2A 3120
L2A 3130
L2A 3140
L2A 3150
L2A 3160
L2A 3170
L2A 3180
L2A 3190
L2A 3200
L2A 3210
L2A 3220
L2A 3230
L2A 3240
L2A 3250
L2A 3260
L2A 3270
L2A 3280
L2A 3290
L2A 3300
L2A 3310
L2A 3320
L2A 3330
L2A 3340
L2A 3350
L2A 3360
L2A 3370
L2A 3380
L2A 3390
L2A 3400
L2A 3410
L2A 3420
L2A 3430
L2A 3440
L2A 3450
L2A 3460
L2A 3470
L2A 3480
L2A 3490
L2A 3500
L2A 3510
L2A 3520
L2A 3530
L2A 3540
L2A 3550
L2A 3560
L2A 3570
L2A 3580
L2A 3590
L2A 3600
L2A 3610
L2A 3620
L2A 3630
L2A 3640
L2A 3650
L2A 3660
L2A 3670
L2A 3680
L2A 3690
L2A 3700

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SUBROUTINE L2B(M, N, M1, L, A, B, W, TOL, MM, NN, MMPNN, L2B 10
* N1, IPIVOT, X, RES, QR, C, IFAULT) L2B 20
C ** PURPOSE ** L2B 30
C SUBROUTINE L2B COMPUTES LEAST SQUARES SOLUTIONS TO OVERDETERMINED L2B 40
C AND UNDERDETERMINED SYSTEMS OF LINEAR EQUATIONS. THE METHOD USED IS L2B 50
C A MODIFIED GRAM-SCHMIDT ORTHOGONAL DECOMPOSITION WITH ITERATIVE L2B 60
C REFINEMENT OF THE SOLUTION. THE SOLUTION MAY BE SUBJECT TO LINEAR L2B 70
C EQUALITY CONSTRAINTS. OUTPUT INCLUDES THE LEAST SQUARES L2B 80
C COEFFICIENTS, RESIDUALS, UNSCALED COVARIANCE MATRIX, AND INFORMATION L2B 90
C ON THE BEHAVIOR OF THE ITERATIVE REFINEMENT PROCEDURE. L2B 100
C MATRIX A IS THE GIVEN MATRIX OF A SYSTEM OF M LINEAR EQUATIONS IN N L2B 110
C UNKNOWN, AND MATRIX W IS A GIVEN DIAGONAL MATRIX OF WEIGHTS WITH ALL L2B 120
C DIAGONAL ELEMENTS NONNEGATIVE. LET H = (SQRT(W))*A. L2B 130
C IN THE EVENT THAT N1 (THE COMPUTED RANK OF MATRIX H) IS LESS THAN N L2B 140
C (THE NUMBER OF UNKNOWN COEFFICIENTS), A UNIQUE SOLUTION VECTOR HAVING L2B 150
C N ELEMENTS CAN BE OBTAINED BY IMPOSING THE CONDITION THAT THE L2B 160
C SOLUTION BE OF MINIMAL EUCLIDEAN NORM. SUCH A SOLUTION IS SOUGHT IN L2B 170
C THE CASE OF UNDERDETERMINED OR RANK-DEFICIENT PROBLEMS. L2B 180
C L2B 190
C ** INPUT VARIABLES ** L2B 200
C M TOTAL NUMBER OF EQUATIONS. L2B 210
C N NUMBER OF UNKNOWN COEFFICIENTS. L2B 220
C M1 NUMBER OF LINEAR CONSTRAINTS (0.LE.M1.LE.M AND M1.LE.N). L2B 230
C L NUMBER OF RIGHT-HAND SIDES (VECTORS OF OBSERVATIONS). L2B 240
C A TWO-DIMENSIONAL ARRAY OF SIZE (MM,N). ON ENTRY, THE ARRAY A L2B 250
C CONTAINS THE GIVEN MATRIX OF A SYSTEM OF M LINEAR EQUATIONS L2B 260
C IN N UNKNOWN, WHERE THE FIRST M1 EQUATIONS ARE TO BE L2B 270
C SATISFIED EXACTLY. A IS LEFT INTACT ON EXIT. L2B 280
C B TWO-DIMENSIONAL ARRAY OF SIZE (MM,L). ON ENTRY, B CONTAINS L2B 290
C THE L GIVEN RIGHT-HAND SIDES (VECTORS OF OBSERVATIONS). B IS L2B 300
C LEFT INTACT ON EXIT. L2B 310
C W VECTOR OF SIZE M. ON ENTRY, W CONTAINS THE DIAGONAL ELEMENTS L2B 320
C OF A GIVEN DIAGONAL MATRIX OF WEIGHTS, ALL NONNEGATIVE. L2B 330
C (THE FIRST M1 ELEMENTS OF W ARE SET EQUAL TO 1.0 BY THE L2B 340
C PROGRAM WHEN M1 IS GREATER THAN ZERO.) ON EXIT, THE ORIGINAL L2B 350
C ELEMENTS OF W HAVE BEEN REPLACED BY THEIR SQUARE ROOTS. L2B 360
C TOL PARAMETER USED IN DETERMINING THE RANK OF MATRIX H. L2B 370
C NOTE -- L2B 380
C (1) IF TOL EQUALS ZERO, THE TOLERANCE USED IN SUBROUTINE L2B 390
C DECOM2 WILL BE BASED ON MACHINE PRECISION. L2B 400
C (2) IF TOL IS GREATER THAN ZERO, THIS VALUE OF TOL WILL BE L2B 410
C USED IN SETTING AN ABSOLUTE TOLERANCE FOR COMPARISON WITH L2B 420
C DIAGONAL ELEMENTS OF THE TRIANGULAR MATRIX OBTAINED IN L2B 430
C SUBROUTINE DECOM2. THE VALUE OF TOL CAN BE BASED ON L2B 440
C KNOWLEDGE CONCERNING THE ACCURACY OF THE DATA. L2B 450
C MM DIMENSIONING PARAMETER SPECIFYING MAXIMUM NUMBER OF ROWS IN L2B 460
C THE ARRAYS A, B AND RES. MM MUST SATISFY MM.GE.M. L2B 470
C NN DIMENSIONING PARAMETER SPECIFYING MAXIMUM NUMBER OF ROWS IN L2B 480
C THE ARRAY X. NN MUST SATISFY NN.GE.N. L2B 490
C MMPNN DIMENSIONING PARAMETER SPECIFYING MAXIMUM NUMBER OF ROWS IN L2B 500
C THE ARRAY QR. MMPNN MUST SATISFY MMPNN.GE.M+N. L2B 510
C L2B 520
C ** OUTPUT VARIABLES AND INTERNAL VARIABLES ** L2B 530
C N1 COMPUTED RANK OF MATRIX H, WHERE H = (SQRT(W))*A. L2B 540
C IPIVOT VECTOR OF SIZE N. ON EXIT, THIS ARRAY RECORDS THE ORDER L2B 550
C IN WHICH THE COLUMNS OF H WERE SELECTED BY THE PIVOTING L2B 560
C SCHEME IN THE COURSE OF THE ORTHOGONAL DECOMPOSITION. L2B 570
C WHENEVER N1.LT.N, THE FIRST N1 ELEMENTS OF IPIVOT INDICATE L2B 580
C WHICH COLUMNS OF H WERE FOUND TO BE LINEARLY INDEPENDENT. L2B 590
C X TWO-DIMENSIONAL ARRAY OF SIZE (NN,L). ON EXIT, X CONTAINS L2B 600
C THE SOLUTION VECTORS. L2B 610
C RES TWO-DIMENSIONAL ARRAY OF SIZE (MM,L). ON EXIT, RES CONTAINS L2B 620
C THE RESIDUAL VECTORS. L2B 630
C QR TWO-DIMENSIONAL ARRAY OF SIZE (MMPNN,N). ON EXIT, QR L2B 640
C CONTAINS THE LOWER TRIANGULAR PORTION OF THE SYMMETRIC L2B 650
C UNSCALED COVARIANCE MATRIX. (THIS ARRAY IS USED INTERNALLY L2B 660
C TO STORE RESULTS FROM SUBROUTINE DECOM2 WHICH ARE L2B 670
C DESTROYED IN COMPUTING THE COVARIANCE MATRIX.) L2B 680
C C VECTOR HAVING AT LEAST 6*(M+N)+2*L ELEMENTS USED (1) FOR L2B 690
C INTERNAL WORK SPACE AND (2) FOR RETURNING INFORMATION ON THE L2B 700
C BEHAVIOR OF THE ITERATIVE REFINEMENT PROCEDURE. L2B 710
C (A) NUMIT IS THE NUMBER OF ITERATIONS CARRIED OUT DURING THE L2B 720
C ITERATIVE REFINEMENT IN ATTEMPTING TO OBTAIN A SOLUTION L2B 730
C FOR THE K-TH RIGHT-HAND SIDE. L2B 740
C ON EXIT, C(K) = +NUMIT IF THE SOLUTION CONVERGED, AND L2B 750
C C(K) = -NUMIT IF THE SOLUTION FAILED TO CONVERGE. L2B 760

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C      (B) DIGITX GIVES AN ESTIMATE OF THE NUMBER OF CORRECT DIGITS L2B 770
C      IN THE INITIAL SOLUTION OF THE COEFFICIENTS FOR THE K-TH L2B 780
C      RIGHT-HAND SIDE. ON EXIT, C(K+L) = DIGITX. L2B 790
C IFAULT FAULT INDICATOR WHICH IS ZERO IF NO ERRORS WERE ENCOUNTERED L2B 800
C AND POSITIVE IF ERRORS WERE DETECTED OR IF EVIDENCE OF SEVERE L2B 810
C ILL-CONDITIONING WAS FOUND. DIAGNOSTIC MESSAGES ARE PRINTED L2B 820
C FROM SUBROUTINE ERROR. IF IFAULT IS SET EQUAL TO 1, 2, 3, 4, L2B 830
C 5, 6 OR 7, EXECUTION IS TERMINATED. EXECUTION CONTINUES WHEN L2B 840
C IFAULT IS SET EQUAL TO 8, 9 OR 10 PROVIDED THAT A SOLUTION L2B 850
C WAS OBTAINED FOR AT LEAST ONE RIGHT-HAND SIDE. THE VALUE OF L2B 860
C IFAULT IS USED TO INDICATE THE FOLLOWING -- L2B 870
C 0 = NO ERRORS ENCOUNTERED. L2B 880
C 1 = BAD INPUT PARAMETER (M, N OR L). L2B 890
C 2 = BAD INPUT PARAMETER (M1). L2B 900
C 3 = BAD DIMENSION. EITHER M.GT.MM, N.GT.NN OR M+N.GT.MMPNN. L2B 910
C 4 = AT LEAST ONE WEIGHT IS NEGATIVE. L2B 920
C 5 = EITHER MATRIX H OR MATRIX OF CONSTRAINTS EQUALS ZERO. L2B 930
C 6 = CONSTRAINTS ARE LINEARLY DEPENDENT. L2B 940
C 7 = ALL SOLUTIONS FAILED TO CONVERGE. L2B 950
C 8 = SOLUTION FAILED TO CONVERGE FOR AT LEAST ONE RIGHT-HAND L2B 960
C SIDE. L2B 970
C 9 = LARGE NUMBER OF ITERATIONS REQUIRED FOR CONVERGENCE. L2B 980
C 10 = ESTIMATED NUMBER OF DIGITS IN INITIAL SOLUTION OF L2B 990
C COEFFICIENTS IS SMALL. L2B 1000
C 11 = DIAGONAL ELEMENT OF COVARIANCE MATRIX WAS COMPUTED TO BE L2B 1010
C NEGATIVE OWING TO ROUNDING ERROR. L2B 1020
C L2B 1030
C ** SUBROUTINES REQUIRED ** L2B 1040
C SUBROUTINE DECOM2 L2B 1050
C USES MODIFIED GRAM-SCHMIDT ALGORITHM WITH PIVOTING TO L2B 1060
C OBTAIN AN ORTHOGONAL DECOMPOSITION OF THE INPUT MATRIX. L2B 1070
C SUBROUTINE SOLVE2 L2B 1080
C COMPUTES COEFFICIENTS AND RESIDUALS. ITERATIVE REFINEMENT IS L2B 1090
C USED TO IMPROVE THE ACCURACY OF THE INITIAL SOLUTION. L2B 1100
C SUBROUTINE SOLVE3 L2B 1110
C CALLED ONLY BY SUBROUTINE SOLVE2. L2B 1120
C SUBROUTINE COVAR L2B 1130
C COMPUTES UNSCALED COVARIANCE MATRIX OF THE COEFFICIENTS. L2B 1140
C SUBROUTINE ERROR L2B 1150
C PRINTS ERROR DIAGNOSTICS WHEN ERRORS ARE DETECTED OR WHEN L2B 1160
C EVIDENCE OF SEVERE ILL-CONDITIONING IS FOUND. L2B 1170
C L2B 1180
C ** STORAGE REQUIREMENTS ** L2B 1190
C THE STORAGE REQUIRED FOR THE DIMENSIONED ARRAYS IN SUBROUTINE L2B IS L2B 1200
C  $M*(2*N + 2*L + 7) + N*(N + L + 7) + 2*L$  L2B 1210
C LOCATIONS. ALL ARRAYS REQUIRED IN SUBROUTINES CALLED BY L2B ARE L2B 1220
C DECLARED HEREIN AND ARE TRANSMITTED ONLY THROUGH PARAMETER LISTS OF L2B 1230
C CALL-SEQUENCES. L2B 1240
C L2B 1250
C ** PRECISION OF ARITHMETIC CALCULATIONS ** L2B 1260
C SINGLE PRECISION ARITHMETIC IS USED FOR ALL CALCULATIONS EXCEPT THE L2B 1270
C DOUBLE PRECISION ACCUMULATION OF INNER PRODUCTS. (THE VARIABLE SUM L2B 1280
C IS DECLARED TO BE DOUBLE PRECISION IN SUBROUTINES DECOM2, SOLVE2, L2B 1290
C SOLVE3 AND COVAR.) IT IS ESSENTIAL FOR THE SUCCESS OF THE ITERATIVE L2B 1300
C REFINEMENT PROCEDURE THAT INNER PRODUCTS BE ACCUMULATED IN DOUBLE L2B 1310
C PRECISION. L2B 1320
C L2B 1330
C ** CONVERSION OF THE PROGRAM TO DOUBLE PRECISION ** L2B 1340
C ***** L2B 1350
C * ON COMPUTERS HAVING SHORT WORD LENGTH (AS THE IBM 360/370) IT MAY * L2B 1360
C * BE DESIRABLE TO PERFORM ALL CALCULATIONS IN DOUBLE PRECISION. IN * L2B 1370
C * THIS CASE, THE ITERATIVE REFINEMENT PRESENTLY INCLUDED IN SOLVE2 * L2B 1380
C * SHOULD BE OMITTED. * L2B 1390
C * TO CONVERT THE PROGRAM TO DOUBLE PRECISION, THE FOLLOWING * L2B 1400
C * APPROACH IS SUGGESTED. * L2B 1410
C * * L2B 1420
C * 1. VARIABLES PRESENTLY DECLARED TO BE REAL SHOULD BE DECLARED * L2B 1430
C * DOUBLE PRECISION. THOSE TYPED INTEGER, DOUBLE PRECISION AND * L2B 1440
C * LOGICAL SHOULD NOT BE CHANGED. * L2B 1450
C * 2. THE USE OF FAIL, NUMIT AND DIGITX SHOULD BE OMITTED. * L2B 1460
C * 3. DESCRIPTION OF VARIABLE C (AT L2B 690-790) SHOULD READ -- * L2B 1470
C * C VECTOR HAVING AT LEAST 6*(M+N) ELEMENTS USED ONLY FOR * L2B 1480
C * INTERNAL WORK SPACE. * L2B 1490
C * 4. THE VALUE OF ETA (AT L2B 1960) SHOULD BE SET SO THAT IT IS THE * L2B 1500
C * SMALLEST POSITIVE DOUBLE PRECISION NUMBER SUCH THAT 1.0 + ETA * L2B 1510
C * IS GREATER THAN 1.0 IN DOUBLE PRECISION ARITHMETIC. * L2B 1520

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          W(I) = SQRT(W(I))
50 CONTINUE
C
C SET PARAMETERS WHICH ALLOCATE VECTOR C TO CONTAIN CERTAIN FINAL
C RESULTS AND ALSO TO BE USED AS WORK SPACE.
C
C K1 IS STARTING POINT FOR NUMIT AND FAIL, OF LENGTH L.
C K2 IS STARTING POINT FOR DIGITX, OF LENGTH L.
C K3 IS STARTING POINT FOR D, OF LENGTH N.
C K4 IS STARTING POINT FOR K-TH COLUMN OF B, OF LENGTH M.
C K5 IS STARTING POINT FOR K-TH COLUMN OF X, OF LENGTH N.
C K6 IS STARTING POINT FOR K-TH COLUMN OF RES, OF LENGTH M.
C K7 IS STARTING POINT FOR WORK SPACE OF LENGTH M.
C K8 IS STARTING POINT FOR WORK SPACE OF LENGTH M.
C K9 IS STARTING POINT FOR WORK SPACE OF LENGTH N.
C K10 IS STARTING POINT FOR WORK SPACE OF LENGTH N.
C K11 IS STARTING POINT FOR WORK SPACE OF LENGTH M + N.
C K12 IS STARTING POINT FOR WORK SPACE OF LENGTH M + N.
C
      K1 = 1
      K2 = K1 + L
      K3 = K2 + L
      K4 = K3 + N
      K5 = K4 + M
      K6 = K5 + N
      K7 = K6 + M
      K8 = K7 + M
      K9 = K8 + M
      K10 = K9 + N
      K11 = K10 + N
      K12 = K11 + M + N
      K = K12 + M + N - 1
C
C MULTIPLY EACH ROW OF MATRIX A (M BY N) BY ITS APPROPRIATE WEIGHT AND
C STORE THE RESULT IN THE FIRST M ROWS OF ARRAY QR. SET ARRAY C AND
C THE LAST N ROWS OF ARRAY QR EQUAL TO ZERO.
C
      DO 60 I=1,K
          C(I) = 0.0
60 CONTINUE
      MP1 = M + 1
      MPN = M + N
      DO 90 J=1,N
          DO 70 I=1,M
              QR(I,J) = A(I,J)*W(I)
70 CONTINUE
          DO 80 I=MP1,MPN
              QR(I,J) = 0.0
80 CONTINUE
90 CONTINUE
C
C OBTAIN AN ORTHOGONAL DECOMPOSITION OF THE MATRIX STORED IN THE FIRST
C M ROWS OF ARRAY QR AND COMPUTE ITS RANK.
C
      CALL DECOM2(M, N, M1, ETA, TOL, QR, C(K3), N1, IPIVOT, SING,
* MMPNN)
C
      IF (.NOT.SING) GO TO 110
      IF (N1.GT.0) GO TO 100
      IFAULT = 5
      CALL ERROR(IFAULT, K, Z)
      RETURN
100 IFAULT = 6
      CALL ERROR(IFAULT, K, Z)
      RETURN
C
C SEEK A SOLUTION (COEFFICIENTS AND RESIDUALS) FOR EACH OF THE L LEAST
C SQUARES PROBLEMS WHOSE RIGHT-HAND SIDES ARE GIVEN IN THE ARRAY B.
C
110 DO 200 K=1,L
C K-TH RIGHT-HAND SIDE.
      K0 = K4 - 1
      DO 120 I=1,M
          K0 = K0 + 1
          C(K0) = B(I,K)
120 CONTINUE

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L2B 2290
L2B 2300
L2B 2310
L2B 2320
L2B 2330
L2B 2340
L2B 2350
L2B 2360
L2B 2370
L2B 2380
L2B 2390
L2B 2400
L2B 2410
L2B 2420
L2B 2430
L2B 2440
L2B 2450
L2B 2460
L2B 2470
L2B 2480
L2B 2490
L2B 2500
L2B 2510
L2B 2520
L2B 2530
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L2B 2670
L2B 2680
L2B 2690
L2B 2700
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L2B 2890
L2B 2900
L2B 2910
L2B 2920
L2B 2930
L2B 2940
L2B 2950
L2B 2960
L2B 2970
L2B 2980
L2B 2990
L2B 3000
L2B 3010
L2B 3020
L2B 3030
L2B 3040

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C
CALL SOLVE2(M, N, M1, A, C(K4), W, N1, IPIVOT, QR, C(K3),
*   ETA, FAIL, NUMIT, DIGITX,
*   C(K5), C(K6), C(K7), C(K8), C(K9), C(K10), C(K11), C(K12),
*   MM, MMPNN)
C
K0 = K5 - 1
DO 130 J=1,N
  K0 = K0 + 1
  X(J,K) = C(K0)
130 CONTINUE
IF (M1.EQ.0) GO TO 150
DO 140 I=1,M1
  RES(I,K) = 0.0
140 CONTINUE
150 M1P1 = M1 + 1
IF (M1P1.GT.M) GO TO 170
K0 = K6 + M1 - 1
DO 160 I=M1P1,M
  K0 = K0 + 1
  RES(I,K) = C(K0)
160 CONTINUE
170 CONTINUE
C
C FOR RIGHT-HAND SIDES WHERE CONVERGENCE OF A SOLUTION IS REPORTED,
C A CHECK IS MADE FOR EVIDENCE OF SEVERE ILL-CONDITIONING. SUCH
C EVIDENCE IS FURNISHED BY LARGE VALUES OF NUMIT (NUMBER OF ITERATIONS
C BEFORE CONVERGENCE WAS OBTAINED) AND SMALL VALUES OF DIGITX
C (ESTIMATE OF THE NUMBER OF CORRECT DIGITS IN THE INITIAL SOLUTION
C OF THE COEFFICIENTS). IF NUMIT EXCEEDS -ALOG10(ETA) A DIAGNOSTIC
C MESSAGE IS PRINTED TO WARN OF ILL-CONDITIONING. IF DIGITX IS LESS
C THAN 0.5 (HALF A DECIMAL DIGIT) A SIMILAR DIAGNOSTIC MESSAGE IS
C PRINTED.
C
C(K) = FLOAT(NUMIT)
IF (FAIL) C(K) = -C(K)
K0 = K2 + K - 1
C(K0) = DIGITX
IF (.NOT.FAIL) GO TO 180
C KSUM IS A TALLY OF SOLUTIONS WHICH FAILED TO CONVERGE.
KSUM = KSUM + 1
IF AULT = 8
CALL ERROR(IFAULT, K, Z)
GO TO 200
180 Z = -ALOG10(ETA)
IF (FLOAT(NUMIT).LE.Z) GO TO 190
IFAULT = 9
Z = FLOAT(NUMIT)
CALL ERROR(IFAULT, K, Z)
190 IF (DIGITX.GE.0.5) GO TO 200
IFAULT = 10
Z = DIGITX
CALL ERROR(IFAULT, K, Z)
200 CONTINUE
IF (KSUM.LT.L) GO TO 210
IFAULT = 7
CALL ERROR(IFAULT, K, Z)
RETURN
210 IF (N1.LT.N) RETURN
DO 230 I=1,N
  MPI = M + I
  DO 220 J=1,N
    QR(I,J) = QR(MPI,J)
220 CONTINUE
QR(I,I) = 0.0
230 CONTINUE
C
C COMPUTE THE UNSCALED COVARIANCE MATRIX OF THE COEFFICIENTS.
C
CALL COVAR(N, M1, N1, IPIVOT, QR, C(K3), C(K9), MMPNN)
C
C IN CERTAIN PROBLEMS, SOME DIAGONAL TERMS OF THE UNSCALED COVARIANCE
C MATRIX ARE EQUAL TO ZERO OR TO SMALL POSITIVE NUMBERS. BECAUSE OF
C ROUNDING ERRORS, COMPUTED VALUES FOR THESE TERMS MAY BE SMALL
C NEGATIVE NUMBERS. AN ERROR DIAGNOSTIC IS PRINTED IF ANY DIAGONAL
C TERM IS NEGATIVE.

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C		L2B 3810
	DO 240 J=1,N	L2B 3820
	IF (QR(J,J).GE.0.0) GO TO 240	L2B 3830
	IFAUULT = 11	L2B 3840
	Z = QR(J,J)	L2B 3850
	CALL ERROR(IFAULT, J, Z)	L2B 3860
240	CONTINUE	L2B 3870
	RETURN	L2B 3880
	END	L2B 3890
C	SUBROUTINE DECOM1(...)	DC1 10
C	SUBROUTINE DECOM1 USES A MODIFIED GRAM-SCHMIDT ALGORITHM WITH	DC1 20
C	PIVOTING TO OBTAIN AN ORTHOGONAL DECOMPOSITION OF THE INPUT MATRIX	DC1 30
C	GIVEN IN Q.	DC1 40
C	THE INPUT PARAMETER TOL (EQUAL EITHER TO ZERO OR TO A POSITIVE	DC1 50
C	NUMBER) IS USED IN DETERMINING THE RANK OF MATRIX Q.	DC1 60
C	NOTE --	DC1 70
C	(1) IF TOL EQUALS ZERO, THE TOLERANCE USED AT STATEMENT DC1 1080	DC1 80
C	WILL BE BASED ON MACHINE PRECISION.	DC1 90
C	UNDER THIS APPROACH, THE TOLERANCE (TOL2) IS SET EQUAL TO	DC1 100
C	(FLOAT(N)*ETA)**2*D(M1+1) AT STATEMENT DC1 1070.	DC1 110
C	IF DESIRED, THE USER CAN OBTAIN A MORE CONSERVATIVE	DC1 120
C	TOLERANCE BY REPLACING N IN THIS STATEMENT BY A LARGER	DC1 130
C	QUANTITY.	DC1 140
C	(2) IF TOL IS GREATER THAN ZERO, TOL2 (EQUAL TO THE SQUARE OF	DC1 150
C	TOL) WILL BE USED AT STATEMENT DC1 1080 AS AN ABSOLUTE	DC1 160
C	TOLERANCE FOR COMPARISON WITH DIAGONAL ELEMENTS OF THE	DC1 170
C	TRIANGULAR MATRIX OBTAINED IN THE DECOMPOSITION. UNDER THIS	DC1 180
C	APPROACH, THE VALUE OF TOL CAN BE BASED ON KNOWLEDGE	DC1 190
C	CONCERNING THE ACCURACY OF THE DATA.	DC1 200
C	ON EXIT, THE ARRAYS Q, R, D AND IPIVOT CONTAIN THE RESULTS OF THE	DC1 210
C	DECOMPOSITION WHICH ARE NEEDED FOR OBTAINING AN INITIAL SOLUTION	DC1 220
C	AND FOR ITERATIVE REFINEMENT.	DC1 230
C	ON EXIT, N1 IS THE COMPUTED RANK OF THE INPUT MATRIX Q.	DC1 240
C	ON EXIT, SING IS SET EQUAL TO .TRUE. WHENEVER	DC1 250
C	(1) N1 = 0 (I.E., INPUT MATRIX Q EQUALS ZERO OR MATRIX OF	DC1 260
C	CONSTRAINTS EQUALS ZERO), OR	DC1 270
C	(2) N1 IS LESS THAN M1 (I.E., THE M1 BY N MATRIX OF LINEAR	DC1 280
C	CONSTRAINTS IS SINGULAR).	DC1 290
C	OTHERWISE, ON EXIT FROM DECOM1, SING = .FALSE.	DC1 300
C	ON EXIT, THE VECTOR IPIVOT RECORDS THE ORDER IN WHICH THE COLUMNS	DC1 310
C	OF Q WERE SELECTED BY THE PIVOTING SCHEME IN THE COURSE OF THE	DC1 320
C	ORTHOGONAL DECOMPOSITION.	DC1 330
	SUBROUTINE DECOM1(M, N, M1, ETA, TOL, Q, R, D, N1, IPIVOT, SING,	DC1 340
	* MM, NN)	DC1 350
	INTEGER IPIVOT(N)	DC1 360
	REAL C, D(1), DM, DS, ETA, Q(MM,N), R(NN,N), RSJ, TOL, TOL2	DC1 370
	DOUBLE PRECISION SUM	DC1 380
	LOGICAL FSUM, SING	DC1 390
	N1 = 0	DC1 400
	SING = .TRUE.	DC1 410
	FSUM = .TRUE.	DC1 420
	MV = 1	DC1 430
	MH = M1	DC1 440
	IF (TOL.GT.0.0) TOL2 = TOL**2	DC1 450
	DO 10 J=1,N	DC1 460
	D(J) = 0.0	DC1 470
	IPIVOT(J) = J	DC1 480
	10 CONTINUE	DC1 490
C	STEP NUMBER NS OF THE DECOMPOSITION.	DC1 500
	DO 210 NS=1,N	DC1 510
	NSM1 = NS - 1	DC1 520
	NSP1 = NS + 1	DC1 530
	IF (NS.EQ.M1+1) GO TO 20	DC1 540
	GO TO 30	DC1 550
20	IF (M1.EQ.M) GO TO 150	DC1 560
	MV = M1 + 1	DC1 570
	MH = M	DC1 580
	FSUM = .TRUE.	DC1 590
C	PIVOT SEARCH.	DC1 600
30	DS = 0.0	DC1 610
	NP = NS	DC1 620
	DO 80 J=NS,N	DC1 630
	IK = IPIVOT(J)	DC1 640
	IF (FSUM) GO TO 40	DC1 650

40	GO TO 60	DC1 660
	SUM = 0.0	DC1 670
	DO 50 L=MV,MH	DC1 680
	SUM = SUM + DBLE(Q(L,IK))*DBLE(Q(L,IK))	DC1 690
50	CONTINUE	DC1 700
	D(J) = SUM	DC1 710
60	IF (DS.LT.D(J)) GO TO 70	DC1 720
	GO TO 80	DC1 730
70	DS = D(J)	DC1 740
	NP = J	DC1 750
80	CONTINUE	DC1 760
C	END PIVOT SEARCH.	DC1 770
	IK = IPIVOT(NP)	DC1 780
	IF (FSUM) DM = DS	DC1 790
	IF (DS.LT.ETA*DM) GO TO 90	DC1 800
	FSUM = .FALSE.	DC1 810
	GO TO 100	DC1 820
90	FSUM = .TRUE.	DC1 830
100	IF (FSUM) GO TO 30	DC1 840
	IF (NP.NE.NS) GO TO 110	DC1 850
	GO TO 130	DC1 860
C	COLUMN INTERCHANGE.	DC1 870
110	IPIVOT(NP) = IPIVOT(NS)	DC1 880
	IPIVOT(NS) = IK	DC1 890
	D(NP) = D(NS)	DC1 900
	IF (NS.EQ.1) GO TO 130	DC1 910
	DO 120 L=1,NSM1	DC1 920
	C = R(L,NP)	DC1 930
	R(L,NP) = R(L,NS)	DC1 940
	R(L,NS) = C	DC1 950
120	CONTINUE	DC1 960
C	END COLUMN INTERCHANGE.	DC1 970
C	RETURN HERE IF N1 = 0. EITHER INPUT MATRIX Q EQUALS ZERO OR MATRIX	DC1 980
C	OF CONSTRAINTS EQUALS ZERO.	DC1 990
130	IF (NS.EQ.1 .AND. DS.EQ.0.0) RETURN	DC1 1000
	SUM = 0.0	DC1 1010
	DO 140 L=MV,MH	DC1 1020
	SUM = SUM + DBLE(Q(L,IK))*DBLE(Q(L,IK))	DC1 1030
140	CONTINUE	DC1 1040
	D(NS) = SUM	DC1 1050
	DS = D(NS)	DC1 1060
	IF (TOL.EQ.0.0) TOL2 = (FLOAT(N)*ETA)**2*D(M1+1)	DC1 1070
	IF (NS.GT.M1 .AND. DS.LE.TOL2) GO TO 150	DC1 1080
	GO TO 160	DC1 1090
150	SING = .FALSE.	DC1 1100
C	RETURN HERE IF N1.LT.N, N1.GT.0 AND N1.GE.M1.	DC1 1110
	RETURN	DC1 1120
C	RETURN HERE IF MATRIX OF CONSTRAINTS IS FOUND TO BE SINGULAR.	DC1 1130
160	IF (DS.EQ.0.0) RETURN	DC1 1140
	IF (NSP1.GT.N) GO TO 200	DC1 1150
C	BEGIN ORTHOGONALIZATION.	DC1 1160
	DO 190 J=NSP1,N	DC1 1170
	NP = IPIVOT(J)	DC1 1180
	SUM = 0.0	DC1 1190
	DO 170 L=MV,MH	DC1 1200
	SUM = SUM + DBLE(Q(L,NP))*DBLE(Q(L,IK))	DC1 1210
170	CONTINUE	DC1 1220
	RSJ = SUM	DC1 1230
	RSJ = RSJ/DS	DC1 1240
	R(NS,J) = RSJ	DC1 1250
	DO 180 L=MV,M	DC1 1260
	Q(L,NP) = Q(L,NP) - RSJ*Q(L,IK)	DC1 1270
180	CONTINUE	DC1 1280
	D(J) = D(J) - DS*RSJ*RSJ	DC1 1290
190	CONTINUE	DC1 1300
C	END ORTHOGONALIZATION.	DC1 1310
200	N1 = N1 + 1	DC1 1320
210	CONTINUE	DC1 1330
C	END STEP NUMBER NS.	DC1 1340
	SING = .FALSE.	DC1 1350
C	NORMAL RETURN. N1 = N.	DC1 1360
	RETURN	DC1 1370
	END	DC1 1380

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C      SUBROUTINE DECOM2(...) DC2 10
C SUBROUTINE DECOM2 USES A MODIFIED GRAM-SCHMIDT ALGORITHM WITH DC2 20
C PIVOTING TO OBTAIN AN ORTHOGONAL DECOMPOSITION OF THE INPUT MATRIX DC2 30
C GIVEN IN QR. DC2 40
C THE INPUT PARAMETER TOL (EQUAL EITHER TO ZERO OR TO A POSITIVE DC2 50
C NUMBER) IS USED IN DETERMINING THE RANK OF MATRIX QR. DC2 60
C NOTE -- DC2 70
C (1) IF TOL EQUALS ZERO, THE TOLERANCE USED AT STATEMENT DC2 1180 DC2 80
C WILL BE BASED ON MACHINE PRECISION. DC2 90
C UNDER THIS APPROACH, THE TOLERANCE (TOL2) IS SET EQUAL TO DC2 100
C (FLOAT(N)*ETA)**2*D(M1+1) AT STATEMENT DC2 1170. DC2 110
C IF DESIRED, THE USER CAN OBTAIN A MORE CONSERVATIVE DC2 120
C TOLERANCE BY REPLACING N IN THIS STATEMENT BY A LARGER DC2 130
C QUANTITY. DC2 140
C (2) IF TOL IS GREATER THAN ZERO, TOL2 (EQUAL TO THE SQUARE OF DC2 150
C TOL) WILL BE USED AT STATEMENT DC2 1180 AS AN ABSOLUTE DC2 160
C TOLERANCE FOR COMPARISON WITH DIAGONAL ELEMENTS OF THE DC2 170
C TRIANGULAR MATRIX OBTAINED IN THE DECOMPOSITION. UNDER THIS DC2 180
C APPROACH, THE VALUE OF TOL CAN BE BASED ON KNOWLEDGE DC2 190
C CONCERNING THE ACCURACY OF THE DATA. DC2 200
C ON EXIT, THE ARRAYS QR, D AND IPIVOT CONTAIN THE RESULTS OF THE DC2 210
C DECOMPOSITION WHICH ARE NEEDED FOR OBTAINING AN INITIAL SOLUTION DC2 220
C AND FOR ITERATIVE REFINEMENT. DC2 230
C ON EXIT, N1 IS THE COMPUTED RANK OF THE INPUT MATRIX QR. DC2 240
C ON EXIT, SING IS SET EQUAL TO .TRUE. WHENEVER DC2 250
C (1) N1 = 0 (I.E., INPUT MATRIX QR EQUALS ZERO OR MATRIX OF DC2 260
C CONSTRAINTS EQUALS ZERO), OR DC2 270
C (2) N1 IS LESS THAN M1 (I.E., THE M1 BY N MATRIX OF LINEAR DC2 280
C CONSTRAINTS IS SINGULAR). DC2 290
C OTHERWISE, ON EXIT FROM DECOM2, SING = .FALSE. DC2 300
C ON EXIT, THE VECTOR IPIVOT RECORDS THE ORDER IN WHICH THE COLUMNS DC2 310
C OF QR WERE SELECTED BY THE PIVOTING SCHEME IN THE COURSE OF THE DC2 320
C ORTHOGONAL DECOMPOSITION. DC2 330
      SUBROUTINE DECOM2(M, N, M1, ETA, TOL, QR, D, N1, IPIVOT, SING, DC2 340
      * MMPNN) DC2 350
      INTEGER IPIVOT(N) DC2 360
      REAL C, D(1), DM, DS, ETA, QR(MMPNN,N), RSJ, TOL, TOL2 DC2 370
      DOUBLE PRECISION SUM DC2 380
      LOGICAL FINIS, FSUM, SING DC2 390
      N1 = 0 DC2 400
      SING = .TRUE. DC2 410
      FSUM = .TRUE. DC2 420
      MV = 1 DC2 430
      MH = M1 DC2 440
      MS = M DC2 450
      MP1 = M + 1 DC2 460
      FINIS = .FALSE. DC2 470
      IF (TOL.GT.0.0) TOL2 = TOL**2 DC2 480
      DO 10 J=1,N DC2 490
        D(J) = 0.0 DC2 500
        IPIVOT(J) = J DC2 510
      10 CONTINUE DC2 520
C STEP NUMBER NS OF THE DECOMPOSITION. DC2 530
      DO 350 NS=1,N DC2 540
        K = M + NS DC2 550
        IF (NS.EQ.M1+1) GO TO 20 DC2 560
        GO TO 30 DC2 570
      20 IF (M1.EQ.M) GO TO 200 DC2 580
        MV = M1 + 1 DC2 590
        MH = M DC2 600
        FSUM = .TRUE. DC2 610
      30 IF (.NOT.FINIS) GO TO 40 DC2 620
        GO TO 150 DC2 630
C PIVOT SEARCH. DC2 640
      40 DS = 0.0 DC2 650
        NP = NS DC2 660
        DO 90 J=NS,N DC2 670
          IF (FSUM) GO TO 50 DC2 680
          GO TO 70 DC2 690
      50 SUM = 0.0 DC2 700
          DO 60 L=MV,MH DC2 710
            SUM = SUM + DBLE(QR(L,J))*DBLE(QR(L,J)) DC2 720
          60 CONTINUE DC2 730
          D(J) = SUM DC2 740
      70 IF (DS.LT.D(J)) GO TO 80 DC2 750
          GO TO 90 DC2 760

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80	DS = D(J)	DC2	770
	NP = J	DC2	780
90	CONTINUE	DC2	790
	IF (FSUM) DM = DS	DC2	800
	IF (DS.LT.ETA*DM) GO TO 100	DC2	810
	FSUM = .FALSE.	DC2	820
	GO TO 110	DC2	830
100	FSUM = .TRUE.	DC2	840
110	IF (FSUM) GO TO 40	DC2	850
	IF (NP.NE.NS) GO TO 120	DC2	860
	GO TO 140	DC2	870
C	COLUMN INTERCHANGE.	DC2	880
120	IK = IPIVOT(NP)	DC2	890
	IPIVOT(NP) = IPIVOT(NS)	DC2	900
	IPIVOT(NS) = IK	DC2	910
	D(NP) = D(NS)	DC2	920
	KM1 = K - 1	DC2	930
	DO 130 L=1,KM1	DC2	940
	C = QR(L,NP)	DC2	950
	QR(L,NP) = QR(L,NS)	DC2	960
	QR(L,NS) = C	DC2	970
130	CONTINUE	DC2	980
C	END COLUMN INTERCHANGE.	DC2	990
C	END PIVOT SEARCH.	DC2	1000
C	RETURN HERE IF N1 = 0. EITHER INPUT MATRIX QR EQUALS ZERO OR	DC2	1010
C	MATRIX OF CONSTRAINTS EQUALS ZERO.	DC2	1020
140	IF (NS.EQ.1 .AND. DS.EQ.0.0) RETURN	DC2	1030
	GO TO 160	DC2	1040
150	MS = K - 1	DC2	1050
	MH = K - 1	DC2	1060
160	IF (FINIS) GO TO 170	DC2	1070
	C = 0.0	DC2	1080
	GO TO 180	DC2	1090
170	C = 1.0	DC2	1100
180	SUM = DBLE(C)	DC2	1110
	DO 190 L=MV,MH	DC2	1120
	SUM = SUM + DBLE(QR(L,NS))*DBLE(QR(L,NS))	DC2	1130
190	CONTINUE	DC2	1140
	D(NS) = SUM	DC2	1150
	DS = D(NS)	DC2	1160
	IF (TOL.EQ.0.0) TOL2 = (FLOAT(N)*ETA)**2*D(M1+1)	DC2	1170
	IF (.NOT.FINIS .AND. NS.GT.M1 .AND. DS.LE.TOL2) GO TO 200	DC2	1180
	GO TO 290	DC2	1190
200	FINIS = .TRUE.	DC2	1200
	MV = M + 1	DC2	1210
	DO 280 NP=NS,N	DC2	1220
	IF (1.GT.M1) GO TO 250	DC2	1230
	DO 210 L=1,M1	DC2	1240
	QR(L,NP) = 0.0	DC2	1250
210	CONTINUE	DC2	1260
	DO 240 J=1,M1	DC2	1270
	SUM = 0.0	DC2	1280
	DO 220 L=1,M	DC2	1290
	SUM = SUM + DBLE(QR(L,J))*DBLE(QR(L,NP))	DC2	1300
220	CONTINUE	DC2	1310
	C = SUM	DC2	1320
	C = C/D(J)	DC2	1330
	DO 230 L=1,M1	DC2	1340
	QR(L,NP) = QR(L,NP) - C*QR(L,J)	DC2	1350
230	CONTINUE	DC2	1360
240	CONTINUE	DC2	1370
250	MPN1 = M + N1	DC2	1380
	DO 270 JJ=MP1,MPN1	DC2	1390
	J = (M + 1) + (M + N1) - JJ	DC2	1400
	SUM = 0.0	DC2	1410
	DO 260 L=J,MPN1	DC2	1420
	LMM = L - M	DC2	1430
	SUM = SUM + DBLE(QR(J,LMM))*DBLE(QR(L,NP))	DC2	1440
260	CONTINUE	DC2	1450
	QR(J,NP) = -SUM	DC2	1460
270	CONTINUE	DC2	1470
280	CONTINUE	DC2	1480
	GO TO 150	DC2	1490
C	RETURN HERE IF MATRIX OF CONSTRAINTS IS FOUND TO BE SINGULAR.	DC2	1500
290	IF (DS.EQ.0.0) RETURN	DC2	1510
	QR(K,NS) = -1.0	DC2	1520

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        NSP1 = NS + 1
        IF (NSP1.GT.N) GO TO 340
C BEGIN ORTHOGONALIZATION.
        DO 330 J=NSP1,N
            SUM = 0.0
            DO 300 L=MV,MH
                SUM = SUM + DBLE(QR(L,J))*DBLE(QR(L,NS))
300    CONTINUE
            RSJ = SUM
            RSJ = RSJ/DS
            QR(K,J) = RSJ
            DO 310 L=1,MS
                QR(L,J) = QR(L,J) - RSJ*QR(L,NS)
310    CONTINUE
            IF (.NOT.FINIS) GO TO 320
            GO TO 330
320    D(J) = D(J) - DS*RSJ*RSJ
330    CONTINUE
C END ORTHOGONALIZATION.
340    IF (.NOT.FINIS) N1 = N1 + 1
350    CONTINUE
C END STEP NUMBER NS.
        SING = .FALSE.
C NORMAL RETURN.
        RETURN
        END
DC2 1530
DC2 1540
DC2 1550
DC2 1560
DC2 1570
DC2 1580
DC2 1590
DC2 1600
DC2 1610
DC2 1620
DC2 1630
DC2 1640
DC2 1650
DC2 1660
DC2 1670
DC2 1680
DC2 1690
DC2 1700
DC2 1710
DC2 1720
DC2 1730
DC2 1740
DC2 1750
DC2 1760
DC2 1770
DC2 1780

C SUBROUTINE SOLVE1(...)
C SUBROUTINE SOLVE1 USES THE ORTHOGONAL DECOMPOSITION STORED IN Q, R,
C D AND IPIVOT TO COMPUTE THE SOLUTION (COEFFICIENTS AND RESIDUALS)
C TO THE LEAST SQUARES PROBLEM WHOSE RIGHT-HAND SIDE IS GIVEN IN B.
C IN THE EVENT THAT N1 (THE COMPUTED RANK OF MATRIX H) IS LESS THAN N
C (THE NUMBER OF UNKNOWN COEFFICIENTS), THE ORIGINAL MATRIX H (M BY N)
C IS REPLACED BY A SMALLER MATRIX (M BY N1) WHOSE COLUMNS ARE LINEARLY
C INDEPENDENT, AND A SOLUTION IS SOUGHT FOR THE SMALLER SYSTEM OF
C EQUATIONS. THUS N - N1 COLUMNS OF THE ORIGINAL MATRIX H ARE DELETED,
C AND COEFFICIENTS CORRESPONDING TO THESE N - N1 DELETED COLUMNS WILL
C BE SET EQUAL TO ZERO.
C IN NORMAL EXITS, THE SOLUTION IS CONTAINED IN THE VECTOR X
C (COEFFICIENTS) AND THE VECTOR RES (RESIDUALS).
C ITERATIVE REFINEMENT IS USED TO IMPROVE THE ACCURACY OF THE INITIAL
C SOLUTION.
C ON EXIT, FAIL IS SET EQUAL TO .TRUE. IF THE SOLUTION FAILS TO
C IMPROVE SUFFICIENTLY. OTHERWISE, FAIL = .FALSE. INFORMATION ON THE
C BEHAVIOR OF THE ITERATIVE REFINEMENT PROCEDURE IS GIVEN BY NUMIT AND
C DIGITX. NUMIT IS THE NUMBER OF ITERATIONS CARRIED OUT IN ATTEMPTING
C TO OBTAIN A SOLUTION. DIGITX IS AN ESTIMATE OF THE NUMBER OF
C CORRECT DIGITS IN THE INITIAL SOLUTION OF THE COEFFICIENTS.
C
C ***** CONVERSION OF THIS SUBROUTINE TO DOUBLE PRECISION *****
C * IF THE PROGRAM IS CONVERTED SO THAT ALL CALCULATIONS ARE DONE IN
C * DOUBLE PRECISION ARITHMETIC, THE ITERATIVE REFINEMENT PRESENTLY
C * INCLUDED IN SOLVE1 SHOULD BE OMITTED, SINCE THE SUCCESS OF THIS
C * PROCEDURE DEPENDS ON COMPUTING INNER PRODUCTS IN GREATER
C * PRECISION THAN OTHER CALCULATIONS.
C * SEE COMMENTS IN SUBROUTINE L2A REGARDING CONVERSION TO DOUBLE
C * PRECISION. IN ADDITION, THE FOLLOWING COMMENTS INDICATE HOW TO
C * OMIT THE ITERATIVE REFINEMENT FROM THIS SUBROUTINE. STATEMENT
C * NUMBERS GIVEN HERE REFER TO THOSE IN THE RIGHT-HAND MARGIN.
C *
C * 1. IN STATEMENT SV1 480 CHANGE REAL TO DOUBLE PRECISION.
C * 2. REPLACE STATEMENT SV1 810 BY A STATEMENT READING
C * 30 DO 50 I=1,M
C * 3. AFTER STATEMENT SV1 1050 INSERT A STATEMENT READING
C * RETURN
C * 4. OMIT STATEMENTS SV1 140-210, 450, 500-510, 530-560, 610,
C * 680-780, 1600-1780 AND 1800-1860.
C *
C *****
C
SUBROUTINE SOLVE1(M, N, M1, A, B, W, N1, IPIVOT, Q, R, D,
* ETA, FAIL, NUMIT, DIGITX,
* X, RES, F, WRES, G, Y, MM, NN)
INTEGER IPIVOT(N)
REAL A(MM,N), B(1), C, D(1), F(1), G(1), Q(MM,N),
SV1 10
SV1 20
SV1 30
SV1 40
SV1 50
SV1 60
SV1 70
SV1 80
SV1 90
SV1 100
SV1 110
SV1 120
SV1 130
SV1 140
SV1 150
SV1 160
SV1 170
SV1 180
SV1 190
SV1 200
SV1 210
SV1 220
SV1 230
SV1 240
SV1 250
SV1 260
SV1 270
SV1 280
SV1 290
SV1 300
SV1 310
SV1 320
SV1 330
SV1 340
SV1 350
SV1 360
SV1 370
SV1 380
SV1 390
SV1 400
SV1 410
SV1 420
SV1 430
SV1 440
SV1 450
SV1 460
SV1 470
SV1 480

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* R(NN,N), RES(1), W(M), WRES(1), X(1), Y(1)          SV1 490
REAL DIGITX, DXNORM, ETA, ETA2, RDR1, RDR2, RDX1, RDX2, RNR, SV1 500
* RNX, XNORM                                           SV1 510
DOUBLE PRECISION SUM                                  SV1 520
LOGICAL FAIL                                         SV1 530
NUMIT = 0                                            SV1 540
KZ = 0                                               SV1 550
ETA2 = ETA*ETA                                       SV1 560
DO 10 I=1,M                                          SV1 570
  F(I) = B(I)*W(I)                                    SV1 580
  WRES(I) = 0.0                                       SV1 590
  RES(I) = 0.0                                       SV1 600
  IF (W(I).EQ.0.0) KZ = KZ + 1                       SV1 610
10 CONTINUE                                          SV1 620
DO 20 J=1,N                                          SV1 630
  X(J) = 0.0                                         SV1 640
  G(J) = 0.0                                         SV1 650
20 CONTINUE                                          SV1 660
K = 0                                               SV1 670
RDX2 = 0.0                                          SV1 680
RDR2 = 0.0                                          SV1 690
C BEGIN K-TH ITERATION STEP.                        SV1 700
30 IF (K.LT.2) GO TO 40                             SV1 710
  IF (((64.*RDX2.LT.RDX1) .AND. (RDX2.GT.ETA2*RNX)) .OR. SV1 720
  * ((64.*RDR2.LT.RDR1) .AND. (RDR2.GT.ETA2*RNR))) GO TO 40 SV1 730
GO TO 300                                           SV1 740
40 RDX1 = RDX2                                       SV1 750
RDR1 = RDR2                                       SV1 760
RDX2 = 0.0                                          SV1 770
RDR2 = 0.0                                          SV1 780
IF (K.EQ.0) GO TO 100                               SV1 790
C NEW RESIDUALS.                                    SV1 800
DO 50 I=1,M                                          SV1 810
  WRES(I) = WRES(I) + F(I)*W(I)                     SV1 820
  IF (W(I).EQ.0.0) GO TO 50                         SV1 830
  RES(I) = RES(I) + F(I)/W(I)                       SV1 840
50 CONTINUE                                          SV1 850
DO 70 NS=1,N1                                       SV1 860
  J = IPIVOT(NS)                                     SV1 870
  X(J) = X(J) + G(NS)                               SV1 880
  SUM = 0.0                                         SV1 890
  DO 60 L=1,M                                       SV1 900
    SUM = SUM + DBLE(A(L,J))*DBLE(WRES(L))          SV1 910
60 CONTINUE                                          SV1 920
G(NS) = -SUM                                        SV1 930
70 CONTINUE                                          SV1 940
DO 90 I=1,M                                          SV1 950
  SUM = 0.0                                         SV1 960
  IF (I.GT.M1) SUM = DBLE(RES(I))                  SV1 970
  DO 80 L=1,N                                       SV1 980
    SUM = SUM + DBLE(A(I,L))*DBLE(X(L))            SV1 990
80 CONTINUE                                          SV1 1000
SUM = SUM - DBLE(B(I))                             SV1 1010
F(I) = -SUM                                         SV1 1020
F(I) = F(I)*W(I)                                    SV1 1030
IF (W(I).EQ.0.0) RES(I) = DBLE(RES(I)) - SUM      SV1 1040
90 CONTINUE                                          SV1 1050
C END NEW RESIDUALS.                                SV1 1060
100 MV = 1                                          SV1 1070
MH = M1                                             SV1 1080
DO 160 NS=1,N1                                       SV1 1090
  J = IPIVOT(NS)                                     SV1 1100
  IF (NS.NE.M1+1) GO TO 110                         SV1 1110
  MV = M1 + 1                                       SV1 1120
  MH = M                                             SV1 1130
110 NSM1 = NS - 1                                    SV1 1140
SUM = -DBLE(G(NS))                                  SV1 1150
IF (1.GT.NSM1) GO TO 130                            SV1 1160
DO 120 L=1,NSM1                                       SV1 1170
  SUM = SUM + DBLE(R(L,NS))*DBLE(Y(L))            SV1 1180
120 CONTINUE                                          SV1 1190
130 Y(NS) = -SUM                                     SV1 1200
C = 0.0                                             SV1 1210
IF (NS.GT.M1) C = -Y(NS)                            SV1 1220
SUM = DBLE(C)                                       SV1 1230
DO 140 L=MV,MH                                       SV1 1240

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	SUM = SUM + DBLE(Q(L,J))*DBLE(F(L))	SV1 1250
140	CONTINUE	SV1 1260
	C = SUM	SV1 1270
	C = C/D(NS)	SV1 1280
	G(NS) = C	SV1 1290
	DO 150 I=MV,M	SV1 1300
	F(I) = F(I) - C*Q(I,J)	SV1 1310
150	CONTINUE	SV1 1320
160	CONTINUE	SV1 1330
	IF (I.GT.M1) GO TO 210	SV1 1340
	DO 170 I=1,M1	SV1 1350
	F(I) = 0.0	SV1 1360
170	CONTINUE	SV1 1370
	DO 200 NS=1,M1	SV1 1380
	J = IPIVOT(NS)	SV1 1390
	SUM = -DBLE(Y(NS))	SV1 1400
	DO 180 L=1,M	SV1 1410
	SUM = SUM + DBLE(Q(L,J))*DBLE(F(L))	SV1 1420
180	CONTINUE	SV1 1430
	C = SUM	SV1 1440
	C = C/D(NS)	SV1 1450
	DO 190 I=1,M1	SV1 1460
	F(I) = F(I) - C*Q(I,J)	SV1 1470
190	CONTINUE	SV1 1480
200	CONTINUE	SV1 1490
210	DO 240 I=1,N1	SV1 1500
	NS = N1 + 1 - I	SV1 1510
	NSP1 = NS + 1	SV1 1520
	SUM = -DBLE(G(NS))	SV1 1530
	IF (NSP1.GT.N1) GO TO 230	SV1 1540
	DO 220 L=NSP1,N1	SV1 1550
	SUM = SUM + DBLE(R(NS,L))*DBLE(G(L))	SV1 1560
220	CONTINUE	SV1 1570
230	G(NS) = -SUM	SV1 1580
240	CONTINUE	SV1 1590
	DO 250 NS=1,N1	SV1 1600
	RDX2 = RDX2 + G(NS)*G(NS)	SV1 1610
250	CONTINUE	SV1 1620
	DO 260 I=1,M	SV1 1630
	RDR2 = RDR2 + F(I)*F(I)	SV1 1640
260	CONTINUE	SV1 1650
	IF (K.NE.0) GO TO 270	SV1 1660
	RNX = RDX2	SV1 1670
	RNR = RDR2	SV1 1680
270	IF (K.NE.1) GO TO 290	SV1 1690
	XNORM = SQRT(RNX)	SV1 1700
	DXNORM = SQRT(RDX2)	SV1 1710
	IF (XNORM.NE.0.0) GO TO 280	SV1 1720
	DIGITX = -ALOG10(ETA)	SV1 1730
	GO TO 290	SV1 1740
280	DIGITX = -ALOG10(AMAX1(DXNORM/XNORM,ETA))	SV1 1750
C	END K-TH ITERATION STEP.	SV1 1760
290	NUMIT = K	SV1 1770
	K = K + 1	SV1 1780
	GO TO 300	SV1 1790
300	IF ((M1+KZ.EQ.M) .AND. (RDX2.GT.4.*ETA2*RNX)) GO TO 310	SV1 1800
	IF ((RDR2.GT.4.*ETA2*RNR) .AND.	SV1 1810
	* (RDX2.GT.4.*ETA2*RNX)) GO TO 310	SV1 1820
	FAIL = .FALSE.	SV1 1830
	RETURN	SV1 1840
310	FAIL = .TRUE.	SV1 1850
	RETURN	SV1 1860
	END	SV1 1870
C	SUBROUTINE SOLVE2(...)	SV2 10
C	SUBROUTINE SOLVE2 USES THE ORTHOGONAL DECOMPOSITION STORED IN QR, D	SV2 20
C	AND IPIVOT TO COMPUTE THE SOLUTION (COEFFICIENTS AND RESIDUALS)	SV2 30
C	TO THE LEAST SQUARES PROBLEM WHOSE RIGHT-HAND SIDE IS GIVEN IN B.	SV2 40
C	IN THE EVENT THAT N1 (THE COMPUTED RANK OF MATRIX H) IS LESS THAN N	SV2 50
C	(THE NUMBER OF UNKNOWN COEFFICIENTS), A UNIQUE SOLUTION VECTOR HAVING	SV2 60
C	N ELEMENTS CAN BE OBTAINED BY IMPOSING THE CONDITION THAT THE	SV2 70
C	SOLUTION BE OF MINIMAL EUCLIDEAN NORM. SUCH A SOLUTION IS SOUGHT IN	SV2 80
C	THE CASE OF UNDERDETERMINED OR RANK-DEFICIENT PROBLEMS.	SV2 90
C	IN NORMAL EXITS, THE SOLUTION IS CONTAINED IN THE VECTOR X	SV2 100
C	(COEFFICIENTS) AND THE VECTOR RES (RESIDUALS).	SV2 110

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C ITERATIVE REFINEMENT IS USED TO IMPROVE THE ACCURACY OF THE INITIAL SV2 120
C SOLUTION. SV2 130
C ON EXIT, FAIL IS SET EQUAL TO .TRUE. IF THE SOLUTION FAILS TO SV2 140
C IMPROVE SUFFICIENTLY. OTHERWISE, FAIL = .FALSE. INFORMATION ON THE SV2 150
C BEHAVIOR OF THE ITERATIVE REFINEMENT PROCEDURE IS GIVEN BY NUMIT AND SV2 160
C DIGITX. NUMIT IS THE NUMBER OF ITERATIONS CARRIED OUT IN ATTEMPTING SV2 170
C TO OBTAIN A SOLUTION. DIGITX IS AN ESTIMATE OF THE NUMBER OF SV2 180
C CORRECT DIGITS IN THE INITIAL SOLUTION OF THE COEFFICIENTS. SV2 190
C THIS SUBROUTINE CALLS SUBROUTINE SOLVE3. SV2 200
C SV2 210
C ***** CONVERSION OF THIS SUBROUTINE TO DOUBLE PRECISION ***** SV2 220
C * IF THE PROGRAM IS CONVERTED SO THAT ALL CALCULATIONS ARE DONE IN * SV2 230
C * DOUBLE PRECISION ARITHMETIC, THE ITERATIVE REFINEMENT PRESENTLY * SV2 240
C * INCLUDED IN SOLVE2 SHOULD BE OMITTED, SINCE THE SUCCESS OF THIS * SV2 250
C * PROCEDURE DEPENDS ON COMPUTING INNER PRODUCTS IN GREATER * SV2 260
C * PRECISION THAN OTHER CALCULATIONS. * SV2 270
C * SEE COMMENTS IN SUBROUTINE L2B REGARDING CONVERSION TO DOUBLE * SV2 280
C * PRECISION. IN ADDITION, THE FOLLOWING COMMENTS INDICATE HOW TO * SV2 290
C * OMIT THE ITERATIVE REFINEMENT FROM THIS SUBROUTINE. STATEMENT * SV2 300
C * NUMBERS GIVEN HERE REFER TO THOSE IN THE RIGHT-HAND MARGIN. * SV2 310
C * * SV2 320
C * 1. IN STATEMENT SV2 470 CHANGE REAL TO DOUBLE PRECISION. * SV2 330
C * 2. REPLACE STATEMENT SV2 880 BY A STATEMENT READING * SV2 340
C * 30 DO 50 I=1,M * SV2 350
C * 3. REPLACE STATEMENTS SV2 1310-1400 BY A STATEMENT READING * SV2 360
C * RETURN * SV2 370
C * 4. OMIT STATEMENTS SV2 120-190, 440, 490-500, 520-550, 650, * SV2 380
C * 750-850, 1650-1830 AND 1850-1910. * SV2 390
C * * SV2 400
C ***** SV2 410
C SUBROUTINE SOLVE2(M, N, M1, A, B, W, N1, IPIVOT, QR, D, SV2 420
* ETA, FAIL, NUMIT, DIGITX, SV2 430
* X, RES, WRES, Y1, Y2, Y, F, G, MM, MMPNN) SV2 440
INTEGER IPIVOT(N) SV2 450
REAL A(MM,N), B(1), C, D(1), F(1), G(1), SV2 460
* QR(MMPNN,N), RES(1), W(M), WRES(1), X(1), Y(1), Y1(1), Y2(1) SV2 470
REAL DIGITX, DXNORM, ETA, ETA2, RDR1, RDR2, RDX1, RDX2, RNR, SV2 480
* RNX, XNORM SV2 490
DOUBLE PRECISION SUM SV2 500
LOGICAL FAIL SV2 510
NUMIT = 0 SV2 520
KZ = 0 SV2 530
ETA2 = ETA*ETA SV2 540
MP1 = M + 1 SV2 550
MPN = M + N SV2 560
NIP1 = N1 + 1 SV2 570
DO 10 I=1,M SV2 580
F(I) = B(I)*W(I) SV2 590
G(I) = 0.0 SV2 600
WRES(I) = 0.0 SV2 610
RES(I) = 0.0 SV2 620
Y1(I) = 0.0 SV2 630
IF (W(I).EQ.0.0) KZ = KZ + 1 SV2 640
10 CONTINUE SV2 650
DO 20 NS=1,N SV2 660
J = M + NS SV2 670
F(J) = 0.0 SV2 680
G(J) = 0.0 SV2 690
X(NS) = 0.0 SV2 700
Y2(NS) = 0.0 SV2 710
20 CONTINUE SV2 720
K = 0 SV2 730
RDX2 = 0.0 SV2 740
RDR2 = 0.0 SV2 750
C BEGIN K-TH ITERATION STEP. SV2 760
30 IF (K.LT.2) GO TO 40 SV2 770
IF (((64.*RDX2.LT.RDX1) .AND. (RDX2.GT.ETA2*RNX)) .OR. SV2 780
* ((64.*RDR2.LT.RDR1) .AND. (RDR2.GT.ETA2*RNR))) GO TO 40 SV2 790
GO TO 270 SV2 800
40 RDX1 = RDX2 SV2 810
RDR1 = RDR2 SV2 820
RDX2 = 0.0 SV2 830
RDR2 = 0.0 SV2 840
IF (K.EQ.0) GO TO 160 SV2 850
C NEW RESIDUALS. SV2 860
SV2 870

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DO 50 I=1,M	SV2 880
WRES(I) = WRES(I) + F(I)*W(I)	SV2 890
IF (W(I).EQ.0.0) GO TO 50	SV2 900
RES(I) = RES(I) + F(I)/W(I)	SV2 910
Y1(I) = Y1(I) + G(I)	SV2 920
50 CONTINUE	SV2 930
DO 100 NS=1,N	SV2 940
J = M + NS	SV2 950
NP = IPIVOT(NS)	SV2 960
X(NP) = X(NP) + F(J)	SV2 970
Y2(NP) = Y2(NP) + G(J)	SV2 980
SUM = -DBLE(X(NP))	SV2 990
DO 60 L=1,M	SV2 1000
SUM = SUM + DBLE(A(L,NP))*DBLE(Y1(L))	SV2 1010
60 CONTINUE	SV2 1020
G(J) = -SUM	SV2 1030
IF (NS.GT.N1) GO TO 70	SV2 1040
GO TO 80	SV2 1050
70 F(J) = 0.0	SV2 1060
GO TO 100	SV2 1070
80 SUM = 0.0	SV2 1080
DO 90 L=1,M	SV2 1090
SUM = SUM + DBLE(A(L,NP))*DBLE(WRES(L))	SV2 1100
90 CONTINUE	SV2 1110
F(J) = -SUM	SV2 1120
100 CONTINUE	SV2 1130
DO 130 I=1,M	SV2 1140
SUM = 0.0	SV2 1150
IF (I.GT.M1) SUM = DBLE(RES(I))	SV2 1160
DO 110 L=1,N	SV2 1170
SUM = SUM + DBLE(A(I,L))*DBLE(X(L))	SV2 1180
110 CONTINUE	SV2 1190
SUM = SUM - DBLE(B(I))	SV2 1200
F(I) = -SUM	SV2 1210
F(I) = F(I)*W(I)	SV2 1220
IF (W(I).EQ.0.0) RES(I) = DBLE(RES(I)) - SUM	SV2 1230
SUM = 0.0	SV2 1240
IF (I.GT.M1) SUM = DBLE(Y1(I))	SV2 1250
DO 120 L=1,N	SV2 1260
SUM = SUM + DBLE(A(I,L))*DBLE(Y2(L))	SV2 1270
120 CONTINUE	SV2 1280
G(I) = -SUM	SV2 1290
130 CONTINUE	SV2 1300
IF (N1P1.GT.N) GO TO 160	SV2 1310
DO 150 I=N1P1,N	SV2 1320
NS = N + N1P1 - I	SV2 1330
J = M + NS	SV2 1340
SUM = 0.0	SV2 1350
DO 140 L=1,J	SV2 1360
SUM = SUM + DBLE(QR(L,NS))*DBLE(G(L))	SV2 1370
140 CONTINUE	SV2 1380
G(J) = SUM	SV2 1390
150 CONTINUE	SV2 1400
C END NEW RESIDUALS.	SV2 1410
C	SV2 1420
160 CALL SOLVE3(F, M1, M, N1, QR, D, Y, MMPNN)	SV2 1430
C	SV2 1440
IF (N1P1.GT.N) GO TO 200	SV2 1450
DO 190 NS=N1P1,N	SV2 1460
J = M + NS	SV2 1470
SUM = DBLE(G(J))	SV2 1480
DO 170 L=MP1,J	SV2 1490
SUM = SUM + DBLE(QR(L,NS))*DBLE(F(L))	SV2 1500
170 CONTINUE	SV2 1510
C = SUM	SV2 1520
C = C/D(NS)	SV2 1530
DO 180 I=1,J	SV2 1540
F(I) = F(I) - C*QR(I,NS)	SV2 1550
180 CONTINUE	SV2 1560
190 CONTINUE	SV2 1570
200 DO 210 J=MP1,MPN	SV2 1580
G(J) = 0.0	SV2 1590
IF (J.LE.M+N1) G(J) = G(J) + F(J)	SV2 1600
210 CONTINUE	SV2 1610
C	SV2 1620
CALL SOLVE3(G, M1, M, N1, QR, D, Y, MMPNN)	SV2 1630

C		SV2 1640
	DO 220 I=1,M	SV2 1650
	RDR2 = RDR2 + F(I)*F(I)	SV2 1660
220	CONTINUE	SV2 1670
	DO 230 I=MP1,MPN	SV2 1680
	RDX2 = RDX2 + F(I)*F(I)	SV2 1690
230	CONTINUE	SV2 1700
	IF (K.NE.0) GO TO 240	SV2 1710
	RNR = RDR2	SV2 1720
	RNX = RDX2	SV2 1730
240	IF (K.NE.1) GO TO 260	SV2 1740
	XNORM = SQRT(RNX)	SV2 1750
	DXNORM = SQRT(RDX2)	SV2 1760
	IF (XNORM.NE.0.0) GO TO 250	SV2 1770
	DIGITX = -ALOG10(ETA)	SV2 1780
	GO TO 260	SV2 1790
250	DIGITX = -ALOG10(AMAX1(DXNORM/XNORM,ETA))	SV2 1800
C	END K-TH ITERATION STEP.	SV2 1810
260	NUMIT = K	SV2 1820
	K = K + 1	SV2 1830
	GO TO 300	SV2 1840
270	IF ((M1+KZ.EQ.M) .AND. (RDX2.GT.4.*ETA2*RNX)) GO TO 280	SV2 1850
	IF ((RDR2.GT.4.*ETA2*RNR) .AND.	SV2 1860
	* (RDX2.GT.4.*ETA2*RNX)) GO TO 280	SV2 1870
	FAIL = .FALSE.	SV2 1880
	RETURN	SV2 1890
280	FAIL = .TRUE.	SV2 1900
	RETURN	SV2 1910
	END	SV2 1920
	SUBROUTINE SOLVE3(F, M1, M, N1, QR, D, Y, MMPNN)	SV3 10
C	SUBROUTINE SOLVE3 IS CALLED ONLY BY SUBROUTINE SOLVE2.	SV3 20
C	THIS SUBROUTINE CALCULATES NEW VALUES OF F.	SV3 30
	REAL C, D(1), F(1), QR(MMPNN,N1), Y(1)	SV3 40
	DOUBLE PRECISION SUM	SV3 50
	MV = 1	SV3 60
	MH = M1	SV3 70
	DO 100 NS=1,N1	SV3 80
	J = M + NS	SV3 90
	IF (NS.EQ.M1+1) GO TO 10	SV3 100
	GO TO 20	SV3 110
10	MV = M1 + 1	SV3 120
	MH = M	SV3 130
20	NSM1 = NS - 1	SV3 140
	SUM = -DBLE(F(J))	SV3 150
	IF (NS.EQ.1) GO TO 40	SV3 160
	DO 30 L=1,NSM1	SV3 170
	MPL = M + L	SV3 180
	SUM = SUM + DBLE(QR(MPL,NS))*DBLE(Y(L))	SV3 190
30	CONTINUE	SV3 200
40	Y(NS) = -SUM	SV3 210
	IF (NS.GT.M1) GO TO 50	SV3 220
	GO TO 60	SV3 230
50	C = -Y(NS)	SV3 240
	GO TO 70	SV3 250
60	C = 0.0	SV3 260
70	SUM = DBLE(C)	SV3 270
	DO 80 L=MV,MH	SV3 280
	SUM = SUM + DBLE(QR(L,NS))*DBLE(F(L))	SV3 290
80	CONTINUE	SV3 300
	C = SUM	SV3 310
	C = C/D(NS)	SV3 320
	F(J) = C	SV3 330
	DO 90 L=MV,M	SV3 340
	F(L) = F(L) - C*QR(L,NS)	SV3 350
90	CONTINUE	SV3 360
100	CONTINUE	SV3 370
	IF (1.GT.M1) GO TO 150	SV3 380
	DO 110 L=1,M1	SV3 390
	F(L) = 0.0	SV3 400
110	CONTINUE	SV3 410
	DO 140 NS=1,M1	SV3 420
	SUM = -DBLE(Y(NS))	SV3 430
	DO 120 L=1,M	SV3 440
	SUM = SUM + DBLE(QR(L,NS))*DBLE(F(L))	SV3 450

120	CONTINUE	SV3	460
	C = SUM	SV3	470
	C = C/D(NS)	SV3	480
	DO 130 L=1,M1	SV3	490
	F(L) = F(L) - C*QR(L,NS)	SV3	500
130	CONTINUE	SV3	510
140	CONTINUE	SV3	520
150	DO 170 NS=1,N1	SV3	530
	J = M + N1 + 1 - NS	SV3	540
	MPN1 = M + N1	SV3	550
	SUM = 0.0	SV3	560
	DO 160 L=J,MPN1	SV3	570
	LMM = L - M	SV3	580
	SUM = SUM + DBLE(QR(J,LMM))*DBLE(F(L))	SV3	590
160	CONTINUE	SV3	600
	F(J) = -SUM	SV3	610
170	CONTINUE	SV3	620
	RETURN	SV3	630
	END	SV3	640
	SUBROUTINE COVAR(N, M1, N1, IPIVOT, C, D, Z, NN)	COV	10
	C SUBROUTINE COVAR USES RESULTS FROM THE ORTHOGONAL DECOMPOSITION	COV	20
	C STORED IN C, D AND IPIVOT TO COMPUTE THE UNSCALED COVARIANCE MATRIX	COV	30
	C OF THE LEAST SQUARES COEFFICIENTS.	COV	40
	C ON ENTRY, THE FIRST N ROWS AND THE FIRST N COLUMNS OF C CONTAIN THE	COV	50
	C UPPER TRIANGULAR MATRIX OBTAINED FROM THE DECOMPOSITION. THIS INPUT	COV	60
	C MATRIX IS DESTROYED IN SUBSEQUENT CALCULATIONS.	COV	70
	C ON EXIT, THE LOWER TRIANGULAR PORTION OF THE SYMMETRIC UNSCALED	COV	80
	C COVARIANCE MATRIX IS CONTAINED IN	COV	90
	C C(1,1)	COV	100
	C C(2,1) C(2,2)	COV	110
	C . . .	COV	120
	C C(N,1) C(N,2) . . . C(N,N)	COV	130
	C IF N1 IS LESS THAN N, ONE OR MORE COLUMNS OF THE MATRIX	COV	140
	C H = (SQRT(W))*A WERE REJECTED AS BEING LINEARLY DEPENDENT. WHENEVER	COV	150
	C THE K-TH COLUMN OF H WAS SO REJECTED, C(I,J) IS SET EQUAL TO ZERO,	COV	160
	C FOR I = K OR J = K, I.GE.J.	COV	170
	INTEGER IPIVOT(N)	COV	180
	REAL C(NN,N), D(1), Z(1)	COV	190
	DOUBLE PRECISION SUM	COV	200
	L = N1	COV	210
	IF (L.GT.M1) C(L,L) = 1.0/D(L)	COV	220
	IF (L.EQ.1) GO TO 60	COV	230
10	J = L - 1	COV	240
	IF (J.GT.M1) C(J,J) = 1.0/D(J)	COV	250
	DO 20 K=L,N1	COV	260
	Z(K) = C(J,K)	COV	270
20	CONTINUE	COV	280
	I = N1	COV	290
	DO 40 KA=J,N1	COV	300
	SUM = 0.0	COV	310
	IF (I.EQ.J) SUM = DBLE(C(I,J))	COV	320
	DO 30 K=L,N1	COV	330
	SUM = SUM - DBLE(Z(K))*DBLE(C(K,I))	COV	340
30	CONTINUE	COV	350
	C(I,J) = SUM	COV	360
	I = I - 1	COV	370
40	CONTINUE	COV	380
	DO 50 K=L,N1	COV	390
	C(J,K) = C(K,J)	COV	400
50	CONTINUE	COV	410
	L = L - 1	COV	420
	IF (L.GT.1) GO TO 10	COV	430
60	IF (N1.EQ.N) GO TO 90	COV	440
	N1P1 = N1 + 1	COV	450
	DO 80 I=1,N	COV	460
	DO 70 J=N1P1,N	COV	470
	C(I,J) = 0.0	COV	480
70	CONTINUE	COV	490
80	CONTINUE	COV	500
	C PERMUTE THE COLUMNS AND ROWS OF MATRIX C TO ACCOUNT FOR PIVOTING.	COV	510
90	DO 120 I=1,N	COV	520
	DO 100 J=1,N	COV	530
	K = IPIVOT(J)	COV	540
	Z(K) = C(I,J)	COV	550

100	CONTINUE	COV	560
	DO 110 J=1,N	COV	570
	C(I,J) = Z(J)	COV	580
110	CONTINUE	COV	590
120	CONTINUE	COV	600
	DO 150 I=1,N	COV	610
	DO 130 J=1,N	COV	620
	K = IPIVOT(J)	COV	630
	Z(K) = C(J,I)	COV	640
130	CONTINUE	COV	650
	DO 140 J=I,N	COV	660
	C(J,I) = Z(J)	COV	670
140	CONTINUE	COV	680
150	CONTINUE	COV	690
	RETURN	COV	700
	END	COV	710
	SUBROUTINE ERROR(IFAULT, K, Z)	ERR	10
C	SUBROUTINE ERROR PRINTS ERROR DIAGNOSTICS IN THE CASE OF ERROR	ERR	20
C	FAILURE.	ERR	30
C	ALSO PRINTED ARE SOME INFORMATIVE DIAGNOSTICS RELATED TO THE	ERR	40
C	ITERATIVE REFINEMENT OF ILL-CONDITIONED SYSTEMS OF EQUATIONS AND TO	ERR	50
C	ROUNDING ERROR PROBLEMS IN COMPUTING THE COVARIANCE MATRIX.	ERR	60
C	IN THE DATA STATEMENT BELOW, NW IS THE PRINTER DEVICE NUMBER.	ERR	70
	REAL Z	ERR	80
	DATA NW /6/	ERR	90
	GO TO (10,20,30,40,50,60,70,80,90,100,110), IFAULT	ERR	100
10	WRITE (NW,99999)	ERR	110
	RETURN	ERR	120
20	WRITE (NW,99998)	ERR	130
	RETURN	ERR	140
30	WRITE (NW,99997)	ERR	150
	IF (K.EQ.2) WRITE (NW,99996)	ERR	160
	RETURN	ERR	170
40	WRITE (NW,99995) K,Z	ERR	180
	RETURN	ERR	190
50	WRITE (NW,99994)	ERR	200
	RETURN	ERR	210
60	WRITE (NW,99993)	ERR	220
	RETURN	ERR	230
70	WRITE (NW,99992)	ERR	240
	RETURN	ERR	250
80	WRITE (NW,99991) K	ERR	260
	RETURN	ERR	270
90	WRITE (NW,99990) K,Z	ERR	280
	RETURN	ERR	290
100	WRITE (NW,99989) K,Z	ERR	300
	RETURN	ERR	310
110	WRITE (NW,99988) K,Z	ERR	320
	RETURN	ERR	330
C	FORMAT STATEMENTS.	ERR	340
99999	FORMAT (50H0*** PARAMETER ERROR. M, N AND L MUST BE GREATER,	ERR	350
	* 11H THAN ZERO.)	ERR	360
99998	FORMAT (50H0*** PARAMETER ERROR. M1 CANNOT EXCEED M OR N, B,	ERR	370
	* 26HUT M1 MUST BE NONNEGATIVE.)	ERR	380
99997	FORMAT (50H0*** DIMENSION ERROR. ONE OR MORE OF THE FOLLOWI,	ERR	390
	* 32HNG ERROR CONDITIONS WAS FOUND --/7X,12HM EXCEEDS MM/7X,	ERR	400
	* 12HN EXCEEDS NN)	ERR	410
99996	FORMAT (5X,17HM+N EXCEEDS MMPNN)	ERR	420
99995	FORMAT (43H0*** WEIGHTS MUST BE NONNEGATIVE. FOR I =,I3,2X,	ERR	430
	* 9HWEIGHT = ,G15.8)	ERR	440
99994	FORMAT (50H0*** EITHER MATRIX H EQUALS ZERO OR MATRIX OF CON,	ERR	450
	* 51HSTRAINTS EQUALS ZERO. NO SOLUTION CAN BE COMPUTED.)	ERR	460
99993	FORMAT (50H0*** SINCE THE CONSTRAINTS ARE LINEARLY DEPENDENT,	ERR	470
	* 29H NO SOLUTION CAN BE COMPUTED.)	ERR	480
99992	FORMAT (39H0*** ALL SOLUTIONS FAILED TO CONVERGE.)	ERR	490
99991	FORMAT (22H0*** FOR B-VECTOR NO.,I3,21H SOLUTION FAILED TO C,	ERR	500
	* 8HONVERGE.)	ERR	510
99990	FORMAT (22H0*** FOR B-VECTOR NO.,I3,21H THE NUMBER OF ITERAT,	ERR	520
	* 45HIONS REQUIRED FOR CONVERGENCE OF SOLUTION WAS,F4.0/4H ***,	ERR	530
	* 57H THIS NUMBER IS LARGE, INDICATING THE PROBLEM IS ILL-CON,	ERR	540
	* 40HDITIONED AND SOLUTION MAY BE INACCURATE.)	ERR	550
99989	FORMAT (22H0*** FOR B-VECTOR NO.,I3,21H ESTIMATED NUMBER OF ,	ERR	560
	* 54HCORRECT DIGITS IN INITIAL SOLUTION OF COEFFICIENTS IS ,	ERR	570
	* G15.8/51H *** SINCE THIS IS SMALL, THE FINAL SOLUTION MAY B,	ERR	580

* 13HE INACCURATE.) ERR 590
 99988 FORMAT (26H0*** DIAGONAL ELEMENT NO.,I3,17H OF THE UNSCALED , ERR 600
 * 57HCOVARIANCE MATRIX WAS COMPUTED TO BE NEGATIVE OWING TO RO, ERR 610
 * 13HUNDING ERROR./29H *** THE COMPUTED VALUE WAS ,G15.8) ERR 620
 END ERR 630

1 MODE 1 *****
 (1) WAMPLER, J.AMER.STAT.ASSN. 1970, P.549, 5TH DEG. POLYNOMIALS, EQUAL WEIGHTS.

21 6 0 2 1 1 1 0.

(F2.0,2F8.0)

0 1 760
 1 6 -2042
 2 63 2111
 3 364 -1684
 4 1365 3888
 5 3906 1858
 6 9331 11379
 7 19608 17560
 8 37449 39287
 9 66430 64382
 10 111111 113159
 11 177156 175108
 12 271453 273291
 13 402234 400186
 14 579195 581243
 15 813616 811568
 16 1118481 1121004
 17 1508598 1506550
 18 2000719 2002767
 19 2613660 2611612
 20 3368421 3369180

1
 (2) FIRST DEGREE POLYNOMIAL, UNEQUAL WEIGHTS.

6 2 0 1 1 2 1 0.

(3F3.0)

1. 2. 2.
 2. 2. 1.
 3. 5. 1.
 4. 4. 1.
 5. 7. 1.
 6. 7. 2.

1
 (3) J. M. CAMERON DATA, UNEQUAL WEIGHTS, TWO COLUMNS LINEARLY DEPENDENT.

7 6 0 2 2 2 1 0.

(2F3.0,F3.1,F3.0,F4.2,F3.0,F5.1,F4.0,F2.0)

1 1 .5 2 .25 2 13.0 130 2
 1 2 .5 2 .25 3 17.0 170 2
 0 3 .0 3 .00 3 18.2 182 1
 0 2 .0 1 .00 1 8.8 88 1
 0 1 .0 -3 .00 0 -3.0 -30 1
 0 1 .0 0 .00 0 2.8 28 1
 0 0 .0 1 .00 0 2.1 21 1

1
 (4) EXAMPLE WITH WEIGHTS AND CONSTRAINTS.

12 6 3 1 2 2 1 0.

(8F3.0)

1 1 1 1 1 1 6 1
 1 1 1 0 0 0 3 1
 1 1 0 0 0 0 2 1
 1 -1 0 0 0 0 1 3
 1 0 -1 0 0 0 -1 3
 1 0 0 -1 0 0 1 3
 1 0 0 0 0 -1 -1 2
 0 1 -1 0 0 0 1 2
 0 1 0 0 -1 0 -1 2
 0 1 0 0 0 -1 1 1
 0 0 1 -1 0 0 -1 1
 0 0 1 0 -1 0 1 1

1
 (5) INVERSE OF HILBERT MATRIX OF ORDER 4. M = 4, N = 4, M1 = 0.

4 4 0 1 2 1 1 0.

(5F7.0)

16. -120. 240. -140. -4.
 -120. 1200. -2700. 1680. 60.
 240. -2700. 6480. -4200. -180.
 -140. 1680. -4200. 2800. 140.

1

(6) INVERSE OF HILBERT MATRIX OF ORDER 4. M = 4, N = 4, M1 = 4.

4 4 4 1 2 1 1 0.
 (5F7.0)
 16. -120. 240. -140. -4.
 -120. 1200. -2700. 1680. 60.
 240. -2700. 6480. -4200. -180.
 -140. 1680. -4200. 2800. 140.
 1

(7) BUSINGER-GOLUB, NUM. MATH. 1965, P.269, INVERSE OF HILBERT MATRIX, ORDER 6.

6 5 1 2 2 1 1 0.
 (5F10.0,10X,2F10.0)
 36. -630. 3360. -7560. 7560. 463. 463.
 -630. 14700. -88200. 211680. -220500. -13860. -17820.
 3360. -88200. 564480. -1411200. 1512000. 97020. 93555.
 -7560. 211680. -1411200. 3628800. -3969000. -258720. -261800.
 7560. -220500. 1512000. -3969000. 4410000. 291060. 288288.
 -2772. 83160. -582120. 1552320. -1746360. -116424. -118944.
 1

(8) EXAMPLE WITH X = 0 (HENCE XNORM = 0). TOL = -1 ON ENTRY TO L2A OR L2B.

1 1 0 1 2 1 0 -1.
 (2F3.0)
 1. 0.
 1

(9) ALBERT, REGRESSION AND THE MOORE-PENROSE INVERSE, 1972, P. 63.

3 4 0 3 2 1 2 0.
 (7F4.0)
 1. 0. 1. 1. 1. 0. 0.
 0. 1. -1. 0. 0. 1. 0.
 1. 1. 0. 1. 0. 0. 1.
 1

(10) FIFTH DEGREE POLYNOMIAL WITH HEAVY WEIGHTS, MATRIX A SCALED. IFAULT=11

21 6 0 1 2 2 1 0.
 (F8.0,3F9.0,2F10.0,F13.2,F10.0)
 1000000. 0. 0. 0. 0. 0. 0. 1000000.1853 16777216.
 1000000. 1000000. 1000000. 1000000. 1000000. 1. -8277497.00 1.
 1000000. 2000000. 4000000. 8000000. 16000000. 32. 85136000.00 1.
 1000000. 3000000. 9000000. 27000000. 81000000. 243. -8245855.00 1.
 1000000. 4000000. 16000000. 64000000. 256000000. 1024. 10500192.00 1.
 1000000. 5000000. 25000000. 125000000. 625000000. 3125. -8191733.00 1.
 1000000. 6000000. 36000000. 216000000. 1296000000. 7776. 8626944.00 1.
 1000000. 7000000. 49000000. 343000000. 2401000000. 16807. -8094491.00 1.
 1000000. 8000000. 64000000. 512000000. 4096000000. 32768. 7897376.00 1.
 1000000. 9000000. 81000000. 729000000. 6561000000. 59049. -7920049.00 1.
 1000000. 10000000. 100000000. 1000000000. 1000000000. 1000000. 6000000.50 16777216.
 1000000. 11000000. 121000000. 1331000000. 14641000000. 161051. -7617047.00 1.
 1000000. 12000000. 144000000. 1728000000. 20736000000. 248832. 8521440.00 1.
 1000000. 13000000. 169000000. 2197000000. 28561000000. 371293. -7113005.00 1.
 1000000. 14000000. 196000000. 2744000000. 38416000000. 537824. 10020992.00 1.
 1000000. 15000000. 225000000. 3375000000. 50625000000. 759375. -6310483.00 1.
 1000000. 16000000. 256000000. 4096000000. 65536000000. 1048576. 12963744.00 1.
 1000000. 17000000. 289000000. 4913000000. 83521000000. 1419857. -5083241.00 1.
 1000000. 18000000. 324000000. 5832000000. 104976000000. 1889568. 12515136.00 1.
 1000000. 19000000. 361000000. 6859000000. 130321000000. 2476099. -3272399.00 1.
 1000000. 20000000. 400000000. 8000000000. 160000000000. 32000000. 6300000.1853 16777216.
 1

(11) LAWSON-HANSON, SOLVING LEAST SQUARES PROBLEMS, 1974, SET 1 EX.16. IFAULT=10

8 6 4 1 2 1 1 0.
 (7F6.0)
 155. 105. -445. -495. -45. -95. -245.
 355. 305. -245. -295. 155. 105. -295.
 -445. -495. -45. -95. 355. 305. 155.
 -245. -295. 155. 105. -445. -495. 105.
 -45. -95. 355. 305. -245. -295. -445.
 155. 105. -445. -495. -45. -95. -495.
 355. 305. -245. -295. 155. 105. -45.
 -445. -495. -45. -95. 355. 305. -95.
 1

(12) LAWSON-HANSON, SOLVING LEAST SQUARES PROBLEMS, P.252, SET 1, EX.16, TOL=.5.

8 6 4 1 2 1 2 0.5
 (7F6.0)
 155. 105. -445. -495. -45. -95. -245.
 355. 305. -245. -295. 155. 105. -295.
 -445. -495. -45. -95. 355. 305. 155.
 -245. -295. 155. 105. -445. -495. 105.
 -45. -95. 355. 305. -245. -295. -445.
 155. 105. -445. -495. -45. -95. -495.

```

355. 305. -245. -295. 155. 105. -45.
-445. -495. -45. -95. 355. 305. -95.
1
(13) BJORCK-GOLUB, BIT 1967, P.322, HILBERT MATRIX INVERSE, ORDER 8. IFAULT=8,9
8 6 2 3 2 1 1 0.
(6F12.0)
20160. -92400. 221760. -288288. 192192. -51480.
945. 945. 8400945.
-952560. 4656960. -11642400. 15567552. -10594584. 2882880.
-40320. -40320. 4159680.
11430720. -58212000. 149688000. -204324120. 141261120. -38918880.
456120. 3256120. 3256120.
-58212000. 304920000. -800415000. 1109908800. -776936160. 216216000.
-2236080. -136080. -136080.
149688000. -800415000. 2134440000. -2996753760. 2118916800. -594594000.
5599440. 7279440. 7279440.
-204324120. 1109908800. -2996753760. 4249941696. -3030051024. 856215360.
-7495488. -6095488. -6095488.
141261120. -776936160. 2118916800. -3030051024. 2175421248. -618377760.
5105100. 6305100. 6305100.
-38918880. 216216000. -594594000. 856215360. -618377760. 176679360.
-1389960. -339960. -339960.
1
(14) LAWSON-HANSON, SOLVING LEAST SQUARES PROBLEMS, 1974, SET 1, EX.12. IFAULT=7
6 8 6 1 2 1 1 0.
(9F6.0)
-245. -295. 155. 105. -445. -495. -45. -95. 355. 1.
355. 305. -245. -295. 155. 105. -445. -495. 305. 1.
-45. -95. 355. 305. -245. -295. 155. 105. -245. 1.
-445. -495. -45. -95. 355. 305. -245. -295. -295. 4.
155. 105. -445. -495. -45. -95. 355. 305. 155. 9.
-245. -295. 155. 105. -445. -495. -45. -95. 105. 16.
1
(15) EXAMPLE WITH SINGULAR MATRIX OF CONSTRAINTS. M1 = 3, N1 = 2. IFAULT=6
6 3 3 1 2 1 1 0.
(4F2.0)
1 1 1 1
2 2 2 1
1 0 0 1
1 2 4 1
1 3 9 1
1 4 9 1
1
(16) EXAMPLE WITH MATRIX A EQUAL TO ZERO (HENCE RANK EQUALS ZERO). IFAULT=5
3 2 0 1 2 1 1 0.
(3F2.0)
0 0 1
0 0 1
0 0 1
1
(17) EXAMPLE WITH ZERO AND NEGATIVE WEIGHTS. IFAULT=4
2 1 0 1 2 2 1 0.
(2F3.0,F4.0)
1. 1. 0.
1. 1. -1.
1
(18) EXAMPLE WHERE N EXCEEDS THE CORRESPONDING DIMENSION LIMIT. IFAULT=3
1 21 0 1 2 1 2 0.
(22F2.0)
1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 1
1
(19) EXAMPLE WHERE M1 EXCEEDS M AND N. IFAULT=2
1 1 2 1 2 1 1 0.
(2F3.0)
1. 1.
1
(20) EXAMPLE WITH M = 0. IFAULT=1
0 1 0 1 2 1 1 0.
(2F3.0)
1. 1. 1.
0
2 MODE 2 *****
(3) J. M. CAMERON DATA, UNEQUAL WEIGHTS, TWO COLUMNS LINEARLY DEPENDENT.
7 6 0 2 2 2 1 0.
(2F3.0,F3.1,F3.0,F4.2,F3.0,F5.1,F4.0,F2.0)
1 1 .5 2 .25 2 13.0 130 2

```

```

1 2 .5 2 .25 3 17.0 170 2
0 3 .0 3 .00 3 18.2 182 1
0 2 .0 1 .00 1 8.8 88 1
0 1 .0 -3 .00 0 -3.0 -30 1
0 1 .0 0 .00 0 2.8 28 1
0 0 .0 1 .00 0 2.1 21 1

```

(9) ALBERT. REGRESSION AND THE MOORE-PENROSE INVERSE, 1972, P. 63.

```

3 4 0 3 2 1 2 0.
(7F4.0)
1. 0. 1. 1. 1. 0. 0.
0. 1. -1. 0. 0. 1. 0.
1. 1. 0. 1. 0. 0. 1.

```

(12) LAWSON-HANSON, SOLVING LEAST SQUARES PROBLEMS, P.252, SET 1, EX.16, TOL=.5.

```

8 6 4 1 2 1 2 0.5
(7F6.0)
155. 105. -445. -495. -45. -95. -245.
355. 305. -245. -295. 155. 105. -295.
-445. -495. -45. -95. 355. 305. 155.
-245. -295. 155. 105. -445. -495. 105.
-45. -95. 355. 305. -245. -295. -445.
155. 105. -445. -495. -45. -95. -495.
355. 305. -245. -295. 155. 105. -45.
-445. -495. -45. -95. 355. 305. -95.

```

(14) LAWSON-HANSON, SOLVING LEAST SQUARES PROBLEMS, 1974, SET 1, EX.12. IFAULT=7

```

6 8 6 1 2 1 1 0.
(9F6.0)
-245. -295. 155. 105. -445. -495. -45. -95. 355. 1.
355. 305. -245. -295. 155. 105. -445. -495. 305. 1.
-45. -95. 355. 305. -245. -295. 155. 105. -245. 1.
-445. -495. -45. -95. 355. 305. -245. -295. -295. 4.
155. 105. -445. -495. -45. -95. 355. 305. 155. 9.
-245. -295. 155. 105. -445. -495. -45. -95. 105. 16.

```

(18) EXAMPLE WHERE N EXCEEDS THE CORRESPONDING DIMENSION LIMIT. IFAULT=3

```

1 21 0 1 2 1 2 0.
(22F2.0)
1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 1
0
1 MODE 1 *****

```

(21) FIRST DEGREE POLYNOMIAL, POSITIVE AND ZERO WEIGHTS. COMPARE EXAMPLE (2).

```

8 2 0 1 1 2 0.
(3F3.0)
1. 2. 2.
2. 2. 1.
3. 5. 1.
4. 4. 1.
5. 7. 1.
6. 7. 2.
7. 9. 0.
8. 6. 0.

```

(22) EXAMPLE WITH WEIGHTS AND CONSTRAINTS. COMPARE EXAMPLE (4).

```

14 6 3 1 2 2 0.
(8F3.0)
1 1 1 1 1 1 6 1
0 0 1 0 -1 0 1 1
1 0 0 0 -1 0 -1 0
0 1 0 -1 0 0 1 0

```

(23) EXAMPLE WITH ZERO WEIGHTS, WHERE RANK A = RANK H = 2. H = (SQRT(W))*A.

```

4 2 0 1 2 2 0.
(4F3.0)
1. 0. 1. 1.
0. 1. 2. 1.
0. 0. 3. 0.
0. 0. 4. 0.

```

(24) EXAMPLE WITH ZERO WEIGHTS, WHERE RANK A = 2, RANK H = 1. H = (SQRT(W))*A.

```

4 2 0 1 2 2 0.
(4F3.0)
1. 0. 1. 1.
0. 1. 2. 0.
0. 0. 3. 0.

```

$\emptyset. \emptyset. 4. 1.$
 \emptyset
1 1 1 \emptyset \emptyset \emptyset 3 1
1 1 \emptyset \emptyset \emptyset \emptyset 2 1
1 -1 \emptyset \emptyset \emptyset \emptyset 1 3
1 \emptyset -1 \emptyset \emptyset \emptyset -1 3
1 \emptyset \emptyset -1 \emptyset \emptyset 1 3
1 \emptyset \emptyset \emptyset \emptyset -1 -1 2
 \emptyset 1 -1 \emptyset \emptyset \emptyset 1 2
 \emptyset 1 \emptyset \emptyset -1 \emptyset -1 2
 \emptyset 1 \emptyset \emptyset \emptyset -1 1 1
 \emptyset \emptyset 1 -1 \emptyset \emptyset -1 1

ALGORITHM 545

An Optimized Mass Storage FFT [C6]

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Key Words and Phrases: multidimensional FFT, fast Fourier transform, FFT, mass storage FFT, mass store sorting, optimal sorting
 CR Categories: 4.9, 5.19, 5.31
 Language: Fortran

DESCRIPTION

The program is an implementation of the optimal sorting algorithm of the author [8] which allows a base-2 version of the Cooley-Tukey FFT algorithm [2-4] efficient access to a mass store array. Optimal sorting for the mass storage FFT has been determined independently by DeLotto and Dotti [5, 6], but in the author's version the emphasis is on "in-place" array modification. This results in slightly higher mass store I/O than the minimum, but requires no additional mass store working space. The method is a logical extension of the work of Singleton [9] and Brenner [1].

The program computes in place the discrete Fourier transform of a one-dimensional or a multidimensional array. In the one-dimensional case the transform is defined by

$$A(J) = \text{SCAL} \sum_{j=0}^{N1-1} a(j) \exp(\pm i 2\pi j J / N1) \quad \text{for } J = 0, 1, \dots, N1 - 1 \quad (1)$$

where SCAL is an arbitrary scaling factor, exp is the exponential function, and $i = \sqrt{-1}$, the sign of the exponent being either plus or minus, depending on the desired transform.

The definition (1) is easily generalized to cover more than one dimension; for example, the two-dimensional case is given by

$$A(K, J) = \text{SCAL} \sum_{k=0}^{N2-1} \sum_{j=0}^{N1-1} a(j, k) \exp(\pm i 2\pi (jJ/N1 + kK/N2))$$

for $K = 0, 1, \dots, N2 - 1$ and $J = 0, 1, \dots, N1 - 1$. (2)

The elements $a(j)$ or $a(j, k)$ in (1) or (2) are the initial complex data in a mass store array. These are replaced by the elements $A(J)$ or $A(K, J)$ as the final complex data. Index reversal in the two-dimensional case is used to indicate dimension transposition by the program (dimension order reversal in general). This may be suppressed in certain cases.

The program consists of a set of subroutines written in ANSI Fortran, only two

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of which are called by a user program. If the data are always complex, the user program calls subroutine CMFFT. If the input data or the result are real, the user program may call the faster subroutine RMFFT. Either subroutine replaces a one-dimensional or a multidimensional array in mass store by its discrete Fourier transform, as defined by eq. (1) or (2).

(If the elements of the mass store array are considered to be singly indexed by $i = 0, 1, 2, \dots$, then the indices of the definitions map into $i = j$ in the one-dimensional case or $i = j + k \cdot N1$ in the two-dimensional case, resultant transposition giving $i = K + J \cdot N2$. In general, $i = j1 + j2 \cdot N1 + j3 \cdot N2 \cdot N1 + \dots$. Remember also that indices are increased by 1 for Fortran.)

The mass store file is assumed to exist and to have been previously defined to the Fortran system and opened for random access by the user. The file is accessed through two system-dependent subroutines MFREAD and MFWRIT (see Appendix F). All other subroutines are system independent.

The user has freedom to specify total array size, mass store block or record size, core store working space size, and the dimensioning of the array, except that all sizes are to base 2 and are given by their binary exponents.

Complex data are transformed by an in-place, base-2 algorithm using postcomputation bit reversal to sort the array [2]. The computation is handled by a modified, in-core FFT routine which does a sequence of partial transforms of the mass store array (the method is discussed further in the following section). The sorting algorithm [8] calculates the most efficient way to access the mass store array for these computations. Finally, the sorting algorithm is used to carry out an overall bit-reversed permutation of the array, again with as few accesses as possible (I/O efficiency is discussed in the conclusion section).

Computation and sorting occur in place, through a combination of "virtual" permutations, where mass store blocks are accessed according to an indexing algorithm but are left physically unpermuted, and symmetric permutations, which interchange blocks according to a generalized index bit reversal. In [8] it is shown that any unsymmetric permutation can be formed from two suitable symmetric permutations, each of which can be done in place.

Multidimensional transforms are achieved automatically by making use of the indexing structure of the FFT algorithm itself. No change to the order of access of elements is necessary, so that the full advantage of the sorting algorithm and program simplicity is maintained.

For transforms in which the initial or final data are real, the usual time-saving algorithm [4] is available to unscramble a half-length complex transform of packed real data (or vice versa). This requires an extra accessing and computing pass through the mass store array but still results in a saving of nearly half the computation and I/O time. In this method it is easy to allow a choice in the degree of redundancy in the final, complex result. The array may be expanded to full redundancy (twice the physical length of the original packed real array), or to partial redundancy, or maintain the same physical length by elimination of all redundancy.

Finally, we must define a number of terms used in the discussion. "Core store" is used to describe a region where the elements of an array are equally accessible at random. "Mass store" implies a region where elements are grouped into "blocks" or "records" which must be accessed as units, but which units are accessible efficiently at random. A "pass" is an array operation which leaves the array elements in place ready for another array operation. Thus an "I/O pass" reads and writes back once all the blocks of an array (this is a logical definition—sometimes in practice not all blocks need be physically accessed, or some may be accessed more than once).

Algorithm Details

General. The structure of the mass storage FFT program is shown in Figure 1. When using real data, routine RMFFT calls CMFFT with a half-length complex

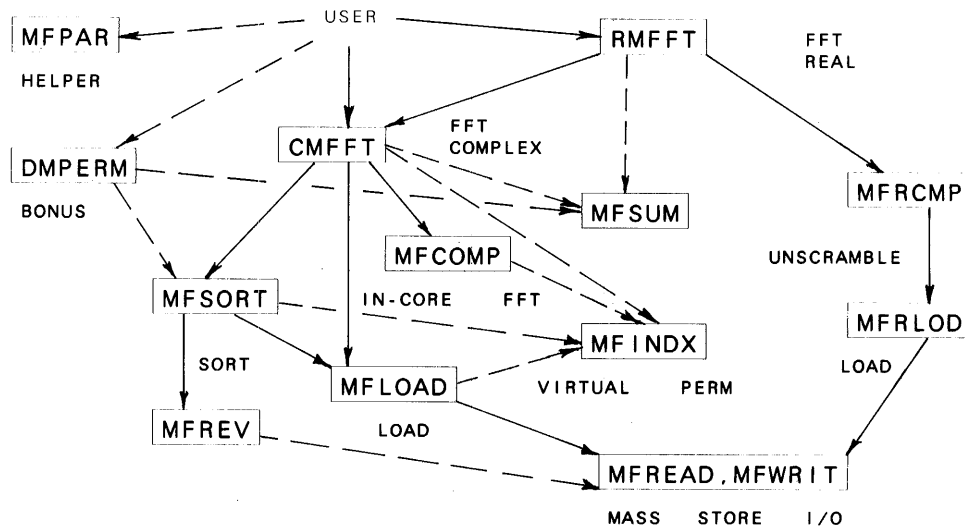


Fig. 1. Internal subroutine calling network. (Solid and dashed lines distinguished roughly between "main path" and "reference" calls.)

array before unscrambling by MFRCMP. If necessary, routines CMFFT and below may be overlaid with routines MFRCMP and MFRLOD. CMFFT controls the sequence of mass store accessing for in-core computation and sorting of the complex FFT. Routines MFCOMP and MFRCMP carry out the FFT computation and unscrambling of data in core store. MFREV, controlled by routine MFSORT, carries out a bit-reversed permutation on elements in core store or on whole blocks in mass store. MFLOAD and MFRLOD load and unload core store for the in-core routines, calling on MFREAD and MFWRIT (Appendix F). MFINDX defines a virtual permutation of the mass store array during the complex FFT operation. MFSUM sums elements in the dimensioning list MEXA(), sorting them if necessary. Routine MFPAR is discussed in Appendix B and routine DMPERM in Appendix D.

The main problem in writing a mass store FFT program is to devise an efficient means of accessing widely spaced data elements for use in the FFT computation kernel and for sorting from bit-reversed order [2]. Singleton [9] describes an algorithm for accessing a mass store array two blocks at a time to carry out each computing pass of an FFT. He organizes mass store accesses in such a way that each I/O and computing pass results in a one-place cyclic shift of array element index bits. This allows the FFT kernel program to access the same elements in core store on each pass, which not only simplifies the program but also brings "widely separated" elements in the original array "within reach" of the in-core computation. Bit-reversed sorting is done in a similar manner. For an array of 2^*M elements, the basic method requires $(2M - 1)$ I/O passes.

With efficient, random access mass store it is possible to use an indexing algorithm to read widely scattered blocks, equivalent to a virtual permutation of the array. After an in-core operation the blocks are written back in place, the blocks remaining physically unpermuted by the operation. In [8] an algorithm is described which uses this method to compute the FFT. The indexing subroutine MFINDX computes the required sequence of block indices using an algorithm which is a generalization of the Block Indexing Algorithm of [8, p. 303]. This is equivalent to a cyclic shift (a number of places) of a set of the block index bits.

Complex Data. Unlike Singleton's algorithm, an attempt is made to load as many blocks together as possible in core store. In this way it is possible to include a number of FFT computing passes while the blocks are in core store, reducing the number of I/O passes by a factor $ICEX - IBEX$ (see the conclusion section). Routine CMFFT uses a slightly modified version of the algorithm of [8, p. 307] for this purpose, doing $M - C$ instead of $M - B$ initial computing passes with

virtual permutations, followed by final C pass computations with direct access, instead of final B pass computations. I/O efficiency is unchanged, but the algorithm is neater (sizes being $2 \times B$ elements per block, $2 \times C$ in core store, $2 \times M$ total, all complex).

Routine MFCOMP computes the FFT in core store. But, unlike other FFT routines, the number of computing passes NPAS, and the effective initial pass number IPAS, are not fixed but are passed by the calling routine CMFFT. Weighting factors are computed using SIN, COS functions on mass store block boundaries but recursively within block boundaries, while array element accessing is ordered to minimize these computations. Multidimensional transforms are achieved by restarting the weighting factor sequence over passes corresponding to the exponent of each dimension, requiring no change to the order of access of elements.

Finally, routine MFSORT is called to sort the array from bit-reversed order. The algorithm is a slightly modified version of the sorting algorithm of [8, pp. 305-307], to allow the reversal of a more general set of index bits. Its structure is very similar to CMFFT, having a first stage using the virtual cyclic shift permutation for block access, with pairs of in-core bit-reversed permutations (MFREV) instead of FFT computations. (The pairs allow unsymmetric cyclic bit shift permutations to be done in place, in core). Only one in-core permutation is required on the first I/O pass. These are followed by a final bit-reversed permutation of whole blocks, if necessary.

In the special case when IDIR = 0, MFSORT is called a number of times to individually reverse bits corresponding to each dimension exponent. This allows the complex FFT to be done without dimension reversal (but note that dimension reversal is required by RMFFT). The dimension shifting subroutine DMPERM, of Appendix D, also uses this feature. In both these cases, I/O is not fully optimized and specially written programs could reduce the number of I/O passes by combining some of the separate, in-core permutations.

Bit-reversed index pairs are generated by a very efficient recursion algorithm (routine MFREV). The algorithm maintains a hierarchy of "reversed" integers of increasing number of bits, up to 2 less than the number being reversed. Incrementing a reversed integer then requires the alternate simple addition of a constant or replacement by the next lower incremented reversed integer in the hierarchy, recursively. For example, with a 3-bit reversal, the reversed set is (0, 4, 2, 6, 1, 5, 3, 7) where the next value is obtained either by adding 4 to the last or by replacing it by one of the values (0, 2, 1, 3) of a 2-bit reversed set.

This method of reversed series generation is in itself fast, as recursion depths are small on the average. But, in addition, only quarter-length series are generated (-2 bits) and the full-length series is derived by scaling by 2 and adding offsets. This is equivalent to reversing an integer $a(\text{integer} - 2 \text{ bits})b$ to $b(\text{reversed integer} - 2 \text{ bits})a$ where a and b are the outer bits. There are four possible combinations for ab , but only those reversals greater numerically than before reversal are required (to prevent nullifying double swaps), leaving in general the three offsets $1 \dots 0$, $0 \dots 0$, and $1 \dots 1$. In particular, only the first offset is required if the internal (reversed integer $- 2$ bits) is smaller than or the same as before reversal. Thus, only valid swap index pairs are generated, saving the unnecessary reversed integer generation of some other methods.

Real Data. For transforms in which the initial or final data is real, a half-length complex transform of packed real data must be unscrambled (or vice versa) by routine MFRCMP. The method relies on complex-conjugate symmetry in the transform of real data [4]. Calculation of indices of multidimensional symmetry is more difficult than in the one-dimensional case. To do this a recursion algorithm is used, which is most easily described by the following Fortran program (for two dimensions, N_1 -by- N_2 elements):

```
L2 = N2/2 + 1
DO 1 J2 = 1, L2
```



```

      K2 = N2 + 2 - J2
      IF(J2.EQ.1)K2 = 1
      L1 = N1
      IF(J2.EQ.K2)L1 = N1/2 + 1
      DO 1 J1 = 1, L1
      K1 = N1 + 2 - J1
      IF(J1.EQ.1)K1 = 1
C   Now have A(K1, K2) and A(J1, J2) as a pair with conjugate symmetry
      1   CONTINUE

```

Note is taken of the special cases which exist when the J index is 1 (Fortran) and when the J and K indices are equal. The general case follows by repetition of the code between the two DO statements, replacing J2 by Jj, etc. In routine MFRCMP arrays JAYA(4), etc., are used for this purpose. To increase the number of dimensions allowed in RMFFT (Appendix A), increase the array sizes. To help visualize the result, the index relationships (Fortran index -1) for an 8-by-4 array are as follows (r for real IF(index(K1, K2).EQ.index(J1, J2)), where index(J1, J2) = J1 + (J2 - 1)*N1, for example, c for complex conjugate):

```

      r00  01  02  03  r04  c03  c02  c01
         10  11  12  13  14  15  16  17
      r20  21  22  23  r24  c23  c22  c21
      c10  c17  c16  c15  c14  c13  c12  c11

```

The array above represents the transformation and transposition of a 4-by-8 real array (imaginary part zero). However, we may also consider it to represent the scrambled transform of an array of initially packed real data, occupying alternately real and imaginary elements, that is, initially 8-by-8 real values. In this case the result is only half-length, element pairs such as 15 and c15 above containing information which can be unscrambled to give a new element 15 and an element 33 in place of c15. In addition, new elements c15 and c33 are derived. Thus the array is expanded to an 8-by-8 full transform of the real data:

```

      r00  01  02  03  r04  c03  c02  c01
         10  11  12  13  14  15  16  17
         20  21  22  23  24  25  26  27
         30  31  32  33  34  35  36  37
      r40  41  42  43  r44  c43  c42  c41
      c30  c37  c36  c35  c34  c33  c32  c31
      c20  c27  c26  c25  c24  c23  c22  c21
      c10  c17  c16  c15  c14  c13  c12  c11

```

Note that the same indexing algorithm applies in both examples (with different N2) but during unscrambling and expansion only the half-length symmetry is used (N2 = 4); corresponding expanded rows in the new array are obtained by a direct offset (4 in this case). Thus, by half-length symmetry, initial rows 1 and 3 are accessed together for unscrambling; these are replaced by unscrambled rows 1 and 3 of the full array while new rows 5 and 7 are also computed and added to the array. It is this ability to add data beyond the existing data which makes dimension reversal essential for in-place computation in the real mass store FFT routine RMFFT.

The complete array is "fully redundant," nearly half the elements being complex conjugates of the other half. If the last three rows of the example are left out, the result is "partially redundant," since row 0 and row 4 still have some internal redundancy. To eliminate all redundancy, rows 0 and 4 can be merged (c43, c42, c41 replacing c03, c02, c01), and pairs of real elements combined as single complex elements (r40 as the imaginary part with r00 real, r44 with r04).

The exact algorithm by which an array with IPAK = 0 or -1 (see Appendix A) can be restored to full redundancy is given by the half-length symmetry program above, putting

```

      A(J1, J2 + N2) = CONJG(A(K1, K2))
      A(K1, K2 + N2) = CONJG(A(J1, J2)), IF(J2.NE.1)

```

with special provision IF(J2.EQ.1) when IPAK = -1:

$$\begin{aligned} A(K1, K2 + N2) &= A(K1, K2) \\ A(J1, J2 + N2) &= \text{CONJG}(A(K1, K2)) \\ A(K1, K2) &= \text{CONJG}(A(J1, J2)) \end{aligned}$$

and IF(index(K1, K2).EQ.index(J1, J2)) when IPAK = -1, values are real:

$$\begin{aligned} A(J1, J2 + N2) &= \text{CMPLX}(\text{AIMAG}(A(J1, J2)), 0.) \\ A(J1, J2) &= \text{CMPLX}(\text{REAL}(A(J1, J2)), 0.) \end{aligned}$$

the general case follows with Jj and Nj instead of J2 and N2 and as before with repetition of the code between DO statements.

Conclusion

The number of I/O passes through the mass store array depends on the array size (2**M), the core store working space size (2**ICEX), and the I/O block or record size (2**IBEX). In general, the larger the working space and the smaller the block size the better, but block size should not be made too small because of other overheads. The FFT computation requires

$$\mathcal{J}((M - \text{IBEX} + 1)/(\text{ICEX} - \text{IBEX}))$$

passes involving I/O (where $\mathcal{J}(x)$ is the smallest integer greater than or equal to x). Postcomputation sorting requires a similar number of passes (without sorting overlap), although DeLotto and Dotti [5] mention

$$\mathcal{J}((M - \text{ICEX} + 1)/(\text{ICEX} - \text{IBEX}))$$

passes. The reason this is not achieved is that in-place sorting requires $M - \text{IBEX} + 1$ virtual permutations to leave a physically unpermuted array. A change in the algorithm to a smaller number $M - \text{ICEX} + 1$ virtual permutations requires an extra block-sorting pass of mass store, nullifying the advantage.

Figure 2 gives the maximum range of block size exponent IBEX necessary to keep the number of I/O passes N of the mass store array low, given ICEX and M. Solid lines bound the optimum $N = 4$. Dashed lines are the boundaries for the next best $N = 6$ passes, and it should not be difficult to operate with $N = 4$ or 6 in most cases. When calling the half-length transform by subroutine RMFFT, increase diagram M and N by 1 to obtain the working M and N.

The hatched line cutting across the figure gives a bound below which all mass store passes access all blocks. In the area above the line,

$$\mathcal{J}((\text{IBEX} - 1)/(\text{ICEX} - \text{IBEX}))$$

sorting passes are required with an additional pass of mass store involving a block reshuffle, in which only some of the blocks are accessed. That is, the optimum can approach $N = 3$ passes. Given an I/O system which allows alteration of the index key of a block, the block shuffle can be replaced by an index shuffle, giving $N - 1$ passes in the upper area. Note also that the diagram is only approximate, as integral rounding effects may increase or decrease the number of passes by one at some points.

Run time depends on two main factors, the computation time of the in-core FFT and the I/O time. A total elapsed time can be written approximately

$$T = \text{TC} * M * 2^{**}M + \text{TM} * N * 2^{**}M$$

where TC and TM are unit computation and unit mass store average access and transfer times corresponding to each complex element. As an example, consider two very different computer systems, a PDP 11/40 system and a CYBER 76 installation. Then we may expect $\text{TC} = 0.5$ ms (PDP 11) or $2 \mu\text{s}$ (CYBER), and $\text{TM} = 1$ ms (PDP 11, $\text{IBEX} = 7$) or 0.5 ms (CYBER, $\text{IBEX} = 9$) for routine CMFFT. Calling subroutine RMFFT with packed real data roughly halves the time, so that a 256-by-256 real array can be transformed in about 5 min (PDP 11) or 1 s (CYBER) CPU time and 8 min (PDP 11) or 1 min (CYBER) elapsed time.

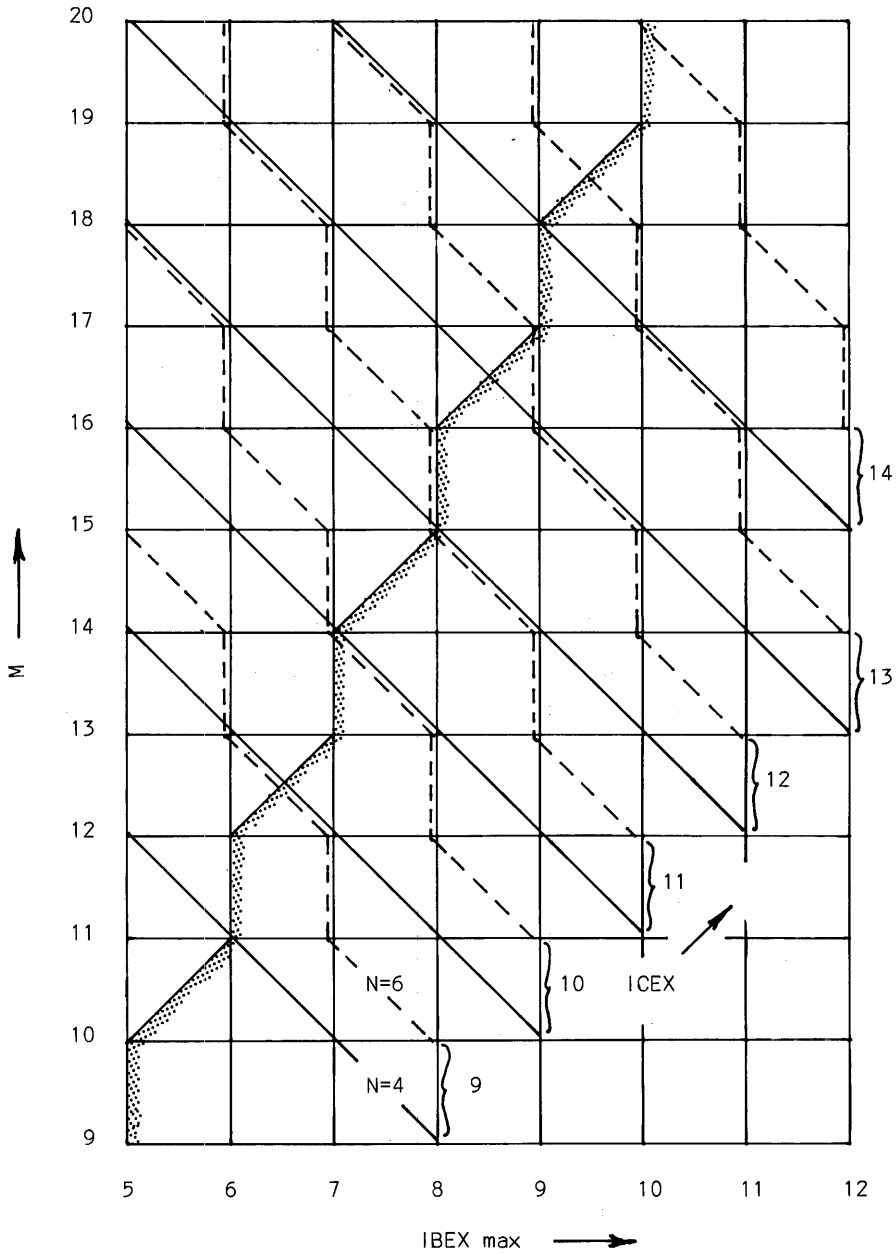


Fig. 2. Mass store I/O passes N related to binary exponents IBEX, ICEx, and M for routine CMFFT. Increase diagram values of M and N by 1 to obtain working M and N for routine RMFFT.

But in machines such as the CYBER, where computing speed is very great compared to I/O latency, the routines are used to great advantage for transforming arrays in LCM (extended memory) in place of mass store, using high-speed, block-copy operations between this and main memory (see Appendixes E and F). Whether an array is multidimensional or not makes little difference to efficiency.

Both TC and TM are block size dependent, there being fixed overheads per block. Both therefore can be written

$$TC = TCO + TCA/2^{**}(IBEX - 1), \quad TM = TMO + TMA/2^{**}(IBEX - 1)$$

where TCA is the FFT computation overhead per block and TMA is the average mass store access time per block. TCO and TMO are the limits for large block size. TCA and TMA set a useful lower limit for block size. In the examples above TM is mainly the result of TMA, or system block access time. Increasing block size reduces TM but with a probable increase in N (depending on M and ICEx).

TC is mainly due to TCO, TCA becoming important typically when $IBEX < 5$. To determine the most efficient parameter combination, begin by choosing ICEX as large as possible; use Figure 2 as a guide in choosing IBEX but follow this up by timing comparisons over a range of IBEX.

The program was written to test the sorting algorithm [8] in practice. At the same time a useful and efficient Fourier program has been obtained, operating with a very general set of parameters. Some parameter combinations can lead to short cuts, not detected or implemented here (see the section, Algorithm Details). Efficiency in a particular case can no doubt be improved by some recoding at the expense of a loss in generality or increase in code complexity.

Hopefully this program will stimulate designers not only of fast transforms but also of general purpose systems. Sorting by index bit manipulation should be considered an important concept for new machine architecture. Index bit reversal is simple to achieve in hardware and, in its generalized form [8], is a symmetric permutation (allowing in-place operation) which forms the basis of a number of useful permutations, not exclusively associated with the FFT.

Appendix A. Program Usage

The complex FFT routine is called by

```
CALL CMFFT (MEXA, NDIM, ISGN, IDIR, SCAL, BUFA, IBEX, ICEX)
```

while the faster, real-to-complex (or vice versa) FFT routine is called by

```
CALL RMFFT (MEXA, NDIM, ISGN, IDIR, SCAL, BUFA, IBEX, ICEX, IPAK)
```

where the subroutine arguments have the following meanings:

MEXA Integer array of size NDIM (defined below). MEXA consists of a list of dimension size binary exponents, defining the dimensioning of the mass store array. For example, a one-dimensional array has a size of $2^{**}MEXA(1)$ elements. A two-dimensional array has $2^{**}MEXA(2)$ sets of $2^{**}MEXA(1)$ adjacent elements each (N1 and N2 of definition equations (1) or (2) are $2^{**}MEXA(1)$ and $2^{**}MEXA(2)$ in initial order).

Notes: (1) The FFT routines actively modify the MEXA list, if necessary, leaving it in the order corresponding to the final array dimensioning. In general, its order is reversed by the FFT, except in the special case with routine CMFFT discussed under parameter IDIR below. The MEXA list should therefore be linked to a unique mass store array so that it always indicates the current dimensioning of that array.

(2) The MEXA list of exponents always refers to the dimensioning of a full array, as required by routine CMFFT or by routine RMFFT with IPAK = +1 (defined below). The complex result of routine RMFFT operating on real data is truncated when IPAK = 0 or -1, so that the MEXA list elements in these cases may refer to a virtual array length.

(3) The first element of the list always gives the size $2^{**}MEXA(1)$ of the set or sets of adjacent elements in the first dimension of the mass store array. If data are of type complex, this is the number of complex elements in each set while if data are of type real (e.g., before calling routine RMFFT in direction real to complex), this is the number of real elements.

(4) Mass store array dimensioning is independent of I/O block transfer size and core working space size.

NDIM Number of dimensions in mass store array (size of MEXA list). Range $1 \leq NDIM \leq 4$ (RMFFT), to increase (see the section, Algorithm Details), or $1 \leq NDIM \leq M$ (CMFFT).

(M) Not a call parameter, but defined here for convenience. $M = \text{sum to NDIM of MEXA list, giving total mass store array size} = 2^{**}M$ elements

- (number of complex elements for routine CMFFT, real or complex elements for routine RMFFT; but see under IPAK packing parameter below).
- ISGN** The sign of ISGN is the sign of the complex exponent of the transform definition (e.g., eq. (1) or (2)). A positive or negative sign results in the inverse transform of the other; but see also below.
- IDIR** Transform direction (RMFFT), reversal (CMFFT) parameter. Routine RMFFT converts packed real data to complex data (or vice versa) during transformation, so that IDIR is needed to determine the direction, independently of ISGN; thus
- IDIR = -1, real to complex (RMFFT), dimension reversal;
IDIR = +1, complex-to-real (RMFFT), dimension reversal.
- (Note that it is usually most convenient to use the same variable for both ISGN and IDIR so that a single negation results in transform inversion.)
- Routine CMFFT does not need a direction parameter, other than ISGN, and IDIR is used in this case to set dimension reversal (transposition in two dimensions) or not, as required, thus
- IDIR not zero, dimension reversal (more efficient I/O);
IDIR = 0, do not use (RMFFT), suppress reversal (CMFFT).
- SCAL** Arbitrary type real scale factor of eq. (1) or (2). If SCAL = 1.0 computation is fastest as no scaling occurs.
- BUFA** Array in core to be used as workspace by FFT routines. Note that internal FFT subroutines assume the following: (1) BUFA is either type real or complex, to suit local needs. (2) BUFA is given the trivial dimension BUFA(1) internally, since its actual size is known only as the exponent ICEX (defined below). (3) Type complex data are assumed to be stored in the sequence real/imaginary/real/imaginary/... in core store and mass store. Some of these points may upset some Fortran compilers.
- IBEX** Mass store I/O block transfer size binary exponent.
Block or record size = 2**IBEX real elements.
- Limits are $2 \leq \text{IBEX} \leq \text{ICEX} - 2$ (RMFFT),
or $1 \leq \text{IBEX} \leq \text{ICEX} - 1$ (CMFFT).
- ICEX** Core store working space size binary exponent.
Dimension of BUFA = 2**ICEX real elements.
- Limits are $\text{IBEX} + 2 \leq \text{ICEX} \leq \text{M}$ (RMFFT),
or $\text{IBEX} + 1 \leq \text{ICEX} \leq \text{M} + 1$ (CMFFT).
- IPAK** Packing parameter (routine RMFFT only). Determines the degree of redundancy (discussed further in the section, Algorithm Details) desired in the complex result after a real-to-complex transform; thus:
- IPAK = +1 gives a fully redundant complex result. The final mass store array is exactly twice the physical length of the initial real array, having the same number, 2**M, of complex elements as initial real elements. This is the same result that is obtained when calling CMFFT with the initial data occupying the real part of a complex array, zero imaginary.
- IPAK = 0 gives a partly redundant complex result, slightly longer than the initial real array. In this case there are $2^{**}(\text{M} - 1) + 2^{**}(\text{M} - \text{MEXA}(\text{NDIM}))$ complex elements in the final array, the MEXA list in the order after the transform (increased to an integral number of mass store blocks, if necessary). This is probably the most useful packing.

IPAK = -1 gives a result containing no redundancy and having exactly the same physical length as the initial 2**M real elements, or 2**(M - 1) complex elements. This is achieved by squeezing together those parts of the array cussed in the section, Algorithm Details. No information is lost and, for one or two dimensions, only a little sorting is needed to access unsqueezed information. For example, in one dimension, the real value at the Nyquist frequency becomes the imaginary part of the zero frequency element, which also must be real.

Notes: (1) For IPAK = +1 or 0 the mass store array file must be extendable. IPAK = -1 has the advantage that the mass store array file remains a fixed length, but has the disadvantage that some user effort is needed to access complex data correctly.

(2) In the complex-to-real direction, IPAK = +1 or 0 are equivalent, since only part of a fully redundant complex array is accessed by RMFFT in this direction. But IPAK = -1 must be used in both directions to correctly handle the squeezed complex array.

Appendix B. Helper Routine

To help the user set up the arguments of Appendix A and to determine the mass store file and block sizes, a helper routine MFPAR is included. Use of the routine is not essential, but is recommended. The routine is called by

```
IERR = MFPAR (IRMF, ICOMP)
```

with IRMF = -1 if mass store data are currently packed real or +1 if data are currently complex when using routine RMFFT. When using routine CMFFT, IRMF = 0. If ICOMP = 0 the argument exponents MEXA(), IBEX, and ICEX are defined by the user while if ICOMP is not zero the exponents are to be computed by MFPAR.

Three COMMON blocks are used to transmit other data; thus

```
COMMON/MFARG/MEXA(4), NDIM, ISGN, IDIR, SCAL, IBEX, ICEX, IPAK
COMMON/MFVAL/DIMA(4), TDM1, RDM1, FBLK, TBLK, RBLK, RCOR, SIZE
COMMON/MFINT/NDMA(4), NTD 1, NRD1, NFBK, NTBK, NRBK, NRCR, NSZE
```

MFARG holds a four-element MEXA() list (for up to four dimensions) followed by the other mass store FFT call arguments (except BUFA). NDIM, IBEX, ICEX, and IPAK must be preset (unless ICOMP not zero), while ISGN, IDIR, and SCAL are ignored here.

COMMON blocks MFVAL and MFINT return data computed by routine MFPAR. These include the four-element arrays DIMA() and NDMA() corresponding element by element to MEXA() but containing the actual dimension sizes 2**MEXA(). Similarly, RBLK, NRBK, RCOR, and NRCR hold the sizes 2**IBEX and 2**ICEX. Note that MFINT variables are one-to-one integer conversions of MFVAL variables, which are of type real. A local variable FIXMAX determines whether a conversion is allowed (set to 32767. by a data statement in MFPAR, but which may be altered to suit). Any values not converted are set to -1 in MFINT and the function returns MFPAR = -1 to indicate this. IBEX and ICEX are forced to be within the limits defined in Appendix A, and MFPAR = +1 if IBEX is forced too small. Otherwise the function returns MFPAR = 0 normally.

If the helper subroutine argument ICOMP is not zero, the computation is reversed and MEXA() exponents are computed from given DIMA() real sizes, IBEX and ICEX computed from RBLK and RCOR real sizes. Sizes are adjusted to be integral powers of 2, adjusted up or down to the closest power on a log scale. NDIM and IPAK must still be given in MFARG.

The most useful values computed by MFPAR are NTD1, NRD1, NFBK, NTBK, and NRBK (in MFINT). NRBK is $2**IBEX$ and is the number of reals in the "record" used for mass store access by the FFT. NTBK is the current total number of records of size NRBK or the current file length. NFBK, on the other hand, is the maximum number of records of size NRBK to be expected, including any mass store array expansion by routine RMFFT (see IPAK = 0 or 1 in Appendix A). NFBK is thus useful for defining the maximum file length to the operating system.

NRD1 is the number of reals in an equivalent record of the current first dimension length ("first" defined by current MEXA(1)). NTD1 is the current number of such equivalent records and these two values are useful for logically accessing a multidimensional mass store array by the user (though not so useful if NDIM = 1). Note that to do this the I/O routines MFREAD and MFWRIT must be able to handle correctly records different in length from NRBK; otherwise NTBK and NRBK should be used. This will be system dependent (see Appendix F).

NSZE, SIZE is the effective total size of the mass store array ($2**M$) and is useful, for example, in computing the scale factor SCAL.

Note that MFPAR should be called just prior to user file access, and with the correct value for IRMF, to ensure the computed variables reflect the current state of the mass store array.

Appendix C. Example

Suppose we wish to transform a two-dimensional real array having 512 rows ($2**9$) of 256 adjacent real elements each ($2**8$), or a total of 128K real elements ($2**17$), where $K = 1024$. Suppose we decide to allow a core store working space of 8K real elements ($2**13$) and to access the mass store array in blocks of 128 real elements ($2**7$) each, these being quite independent of array dimensioning. Then, using the helper routine MFPAR (Appendix B) to compute file parameters, the mass store file is defined and opened for random access, loaded with data, and subroutine RMFFT is called with

```

COMMON/MFARG/MEXA(4), NDIM, ISGN, IDIR, SCAL, IBEX, ICEX, IPAK
COMMON/MFVAL/DIMA(4), TDM1, RDM1, FBLK, TBLK, RBLK, RCOR, SIZE
COMMON/MFINT/NDMA(4), NTD1, NRD1, NFBK, NTBK, NRBK, NRCR, NSZE
C COMMON BLOCKS USED BY HELPER ROUTN MFPAR (NOT ESSENTIAL)
C
  REAL BUFA(8192)
  COMPLEX CBUFA(4096)
  EQUIVALENCE (BUFA(1), CBUFA(1))
C WORKING ARRAY IN CORE (EQUIVALENCED FOR USER ACCESS)
C
  MEXA(1) = 8
  MEXA(2) = 9
  NDIM = 2
  ISGN = -1
  IDIR = -1
  SCAL = 1.0
  IBEX = 7
  ICEX = 13
  IPAK = 1
C MASS STORE FFT ARGUMENTS INITIALIZED
C
  IRMF = - 1
  ICOMP = 0
  IERR = MFPAR (IRMF, ICOMP)
C HELPER ROUTINE, DATA REAL, RMFFT, EXPONENTS DEFINED,
C IF(IERR.EQ.0) OK
C COMPUTES FILE PARAMETERS IN COMMON AREAS (NOT ESSENTIAL)
(open mass store file here)
C HERE CAN OPEN MASS STORE FILE, NFBK MAX 'RECDS' OF NRBK REALS
C

```

```

DO 2 JB = 1, NTD1                (or NTBK)
DO 1 I = 1, NRD1                (or NRBK)
1  BUFA(I) = (Enter a real variable here)
2  CALL MFWRIT (BUFA, NRD1, JB)  (or (BUFA, NRBK, JB))
C  LOAD MASS STORE FILE WITH PACKED REAL DATA
C  (IF MFREAD/MFWRIT REQUIRE FIXED RECD LENGTH, USE NTBK, NRBK)
C

```

CALL RMFFT (MEXA, NDIM, ISGN, IDIR, SCAL, BUFA, IBEX, ICEX, IPAK) resulting in the transform of eq. (2), with SCAL = 1.0, N1 = 256, N2 = 512, and a negative complex exponent. Because the initial data are packed real and because the parameter IPAK = 1 is chosen, the mass store array will be extended to 128K complex elements, or twice its initial physical size. This array is a fully redundant transform of the original, having 256 rows (2**8) of 512 adjacent complex elements (2**9), the dimensions being reversed or transposed. To indicate transposition to the MEXA list will be reversed (MEXA(1) = 9, MEXA(2) = 8 after the transform).

To save unnecessary array extension, the last parameter IPAK can be made 0 or -1. If IPAK = 0 the result is a partially redundant transform, being the first 129 rows (2**MEXA(2)/2 + 1) of the transform, of 512 adjacent complex elements each. The first and last rows each have an internal conjugate symmetry, but other redundancy is deleted.

If IPAK = -1 the first and last rows above are "squeezed" together, the second half of the last row becoming the second half of the first row. In addition, the first and middle real elements of the last row become the imaginary parts of the first and middle elements of the first row, resulting in an array of 128 rows of 512 complex elements, exactly the same physical length as the initial array, so that no file extension is necessary (see the section, Algorithm Details).

The complex result is accessed, then an inverse transform is called by

```

IRMF = 1
IERR = MFPAR (IRMF, ICOMP)
C  HELPER ROUTINE AGAIN, DATA COMPLEX, ROUTINE RMFFT,
C  IF(IERR.EQ.0) OK
C
NCNT = NRD1/2                    (or NRBK/2)
C  NCNT IS THE NUMBER OF COMPLEX ELEMENTS IN RECORD
DO 4 JB = 1, NTD1                (or NTBK)
CALL MFREAD (BUFA, NRD1, JB)    (or (BUFA, NRBK, JB))
DO 3 I = 1, NCNT
3  (access each complex element CBUF(I) here)
4  CALL MFWRIT (BUFA, NRD1, JB)  (or (BUFA, NRBK, JB))
C  COMPLEX RESULT READ BY USER (WRITING NEW VALUES IF DESIRED)
C  (IF MFREAD/MFWRIT REQUIRE FIXED RECD LENGTH, USE NTBK, NRBK)
C
ISGN = -ISGN
IDIR = -IDIR
SCAL = 1.0/SIZE
CALL RMFFT (MEXA, NDIM, ISGN, IDIR, SCAL, BUFA, IBEX, ICEX, IPAK)

```

The scale factor SCAL, using SIZE = 2**M computed by MFPAR, is chosen here to normalize the result to the same scale as the original data. The result is a real array having 512 rows of 256 adjacent real elements each, the MEXA list being restored to its initial order.

Appendix D. Bonus Routine

The sorting algorithm used in the mass storage FFT may also be used for changing the order of dimensioning of a multidimensional mass store array. This is similar to, but more general than, Eklundh's method [7]. As an example, and in addition to the FFT routines, a dimension-shifting routine DMPERM is included for completeness, called by

```
CALL DMPERM (MEXA, NDIM, NSHFT, IREX, BUFA, IBEX, ICEX)
```


where the arguments MEXA, NDIM, BUFA, IBEX, ICEX have the same meaning as in the FFT calls.

The other parameters are

NSHFT Dimension shift count.
 NSHFT = 0 No shift or change occurs.
 NSHFT = 1, 2 etc. First-to-next dimension, circular NSHFT place shift, modulo (NDIM).
 NSHFT = -1 Dimension order reversed.
 IREX "Element size" binary exponent (size = 2**IREX reals); thus IREX = 0 for real array, IREX = 1 for complex array.

Most of the comments concerning the FFT routine parameters apply also to this routine. The mass store array is either real or complex (or "elements" of multiple reals), with an MEXA list defining the current dimensioning (binary exponents). The MEXA list is actively modified by the routine if necessary. (Note that to sort arrays of elements smaller in size than type real, such as the type byte elements of some systems, it is necessary to alter type statements for BUFA and TEMP variables in all subroutines called.)

The routine is quite trivial in design, consisting of only 20 statements. It operates by calling the generalized mass store bit-reversed sorting routine MFSORT, used internally by the FFT routines, and therefore has a similar I/O efficiency to the FFT routines (see the conclusion section).

Appendix E. Test Programs

A universal test program is provided which sets up, for a given M, exhaustive permutations and combinations of the different parameters IBEX, ICEX, and NDIM from one to three dimensions. Mass store is simulated, through dummy routines MFREAD and MFWRIT, by a core store array so that the program is independent of system I/O.

The program stores a pseudorandom number sequence in the simulated mass store, transforms this by the mass store FFT routines, and compares the result with a discrete Fourier transform computed naively by routine NAIVE. An inverse mass store FFT is called and the result compared with the initial data. Maximum differences are noted and various levels of data and difference print-out are available through parameter IPRINT. The test is considered successful if maximum differences are explainable in terms of machine roundoff errors.

Because the tests are exhaustive and because of the naive DFT computation the program is quite slow when the overall array size exponent M is other than very small. For example, in a CYBER 76 computer with M = 5 the program takes 0.3 s while with M = 10 the program takes about 800 s. Less exhaustive testing can be done by replacing program DO-loop variables by fixed variables (see program comments).

Example test programs are also included for accessing mass store in two specific computer systems. These are a CYBER 76 (true mass store and LCM "static" mass store) and a PDP 11. Mass store I/O routines for other systems should be easily devised based on these examples and the comments of Appendix F.

Appendix F. System-Dependent I/O Routines

The system-dependent random access transfers between mass store and core store are handled by two subroutines called internally by the FFT routines; thus

```
CALL MFREAD (BUFA, NB, JB)           read block into core store
CALL MFWRIT (BUFA, NB, JB)         write block from core store
```

where BUFA is the core store address for the start of the transfer, NB is the number of real elements to be transferred, and JB is the desired block or record number. The range of JB is as follows:

for CMFFT or RMFFT if IPAK = +1,

$$1 \leq JB \leq 2^{**}(M - IBEX + 1),$$

for RMFFT if IPAK = 0,

$$1 \leq JB \leq 2^{**}(M - IBEX) + 2^{**}(M - MEXA(NDIM) - IBEX + 1),$$

for MEXA(NDIM) + IBEX < M + 1,

or

$$1 \leq JB \leq 2^{**}(M - IBEX) + 1, \quad \text{for MEXA(NDIM) + IBEX} \geq M + 1,$$

(the MEXA list in the order after the transform)

for RMFFT if IPAK = -1,

$$1 \leq JB \leq 2^{**}(M - IBEX).$$

In addition, the I/O subroutines must know what file to access and any starting block offset, which may be given through COMMON variables, for example. The routines expect the mass store file to be defined and opened. The user program may, of course, call MFREAD and MFWRIT itself to load mass store with data or read the transform result.

For example, in PDP 11 Fortran a random access occurs with

```

SUBROUTINE MFREAD (BUFA, NB, JB)
C  READ BLOCK, INDEX JB, FROM MASS STORE TO BUFA, NB REAL VALUES
  REAL BUFA(NB)
  COMMON/FFTCOM/LUN
  READ (LUN'JB)(BUFA(I), I = 1, NB)  or  READ (LUN'JB) BUFA
  RETURN
END

```

and similarly for subroutine MFWRIT.

As discussed in Appendixes B and C it may be useful to call MFREAD/MFWRIT with a record length NB different from 2**IBEX. This can give a user more logical access to a multidimensional array, in records of first dimension length, for example. Some systems allow a redefinition of the file structure, in which case the problem is solved. Others allow a "word-addressable" file structure in which an I/O transfer may start at any word in the file. In this case, MFREAD/MFWRIT can include the following:

$$INDEX = (JB - 1) * NB + 1$$

(transfer between file real element indices (INDEX) and (INDEX + NB - 1). The dummy I/O routines of the universal test program of Appendix E and the CYBER Extended Core/LCM and PDP 11 Macro routines include this feature.

Subroutines MFREAD and MFWRIT allow the user considerable scope for modifying the operation of the FFT routines. In the universal test driver program (Appendix E) the subroutines simply copy data from one core store array to another. In some systems external core store can be accessed efficiently in blocks, which allows this to be used as a fast, static mass store. Using different files on the first I/O pass for MFREAD and MFWRIT allows data to be copied automatically from a source file to a working and result file. In this case care must be taken to allow only a single read of each block of the source file, any subsequent read, even on the "first" I/O pass, being from the working file (important during complex-to-real transformation by routine RMFFT).

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ALGORITHM

```

C
C   THIS SET OF PROGRAM AND SUBROUTINE UNITS IS TO SUPPORT
C
C   'AN OPTIMIZED MASS STORAGE FFT', BY DONALD FRASER
C
C   REVISION DATE: JULY 1978 (MINOR REV JUNE 79).
C
C   THE SET INCLUDES A UNIVERSAL TEST DRIVER PROGRAM, WHICH SIMULATES
C MASS STORE THROUGH A FORTRAN ARRAY, SAMPLE PROGRAMS AND I/O
C SUBROUTINES FOR CONTROL DATA 6000 AND CYBER COMPUTERS AND FOR
C DEC PDP 11 MINICOMPUTERS.
C
C   I/O SUBROUTINES FOR OTHER SYSTEMS SHOULD BE EASILY CONSTRUCTED
C FROM THE EXAMPLES AND WITH REFERENCE TO THE FORMAL PAPER.
C BUT CARE SHOULD BE TAKEN WITH SOME FUSSY COMPILERS SINCE FFT
C SUBROUTINES ASSIGN EITHER TYPE REAL OR TYPE COMPLEX TO THE SAME
C ARRAY AS IT IS PASSED AS A FORMAL PARAMETER. NOTE THAT COMPLEX
C DATA IS ASSUMED TO BE STORED REAL/IMAG/REAL/IMAG... IN BOTH
C MASS STORE AND CORE STORE.
C
C   THE PROGRAM UNITS APPEAR IN THE FOLLOWING ORDER:
C
C   FIRST, THE FFT SUBROUTINE SET:
C 1 RMFFT      OPTIMIZED MASS STORAGE FFT (REAL DATA OR REAL RESULT)
C 2 CMFFT      CALLED BY 1, OR MASS STORAGE FFT (ALL COMPLEX DATA)
C 3 MFCOMP     IN-CORE FFT
C 4 MFSORT     MASS STORE SORTING
C 5 MFREV     IN-CORE SORTING OR WHOLE BLOCK SORTING
C 6 MFLoad    LOADING/UNLOADING CORE STORE
C 7 MFINDX    BLOCK INDEXING ALGORITHM (VIRTUAL PERMUTATION)
C 8 MFSUM     MEXA() EXPONENT SUMMATIONS
C 9 MFCMP     REAL-COMPLEX UNSCRAMBLING/SCRAMBLING
C 10 MFLod    LOADING/UNLOADING CORE STORE FOR MFCMP
C 11 MFPAR    HELPER ROUTINE (NOT ESSENTIAL, BUT RECOMMENDED)
C 12 DMPerm   MASS STORE DIMENSION SHIFTING (BONUS SUBROUTINE)
C
C   THEN, TEST PROGRAMS AND SAMPLE I/O SUBROUTINES
C
C 13 UNIVERSAL TEST PROGRAM (SIMULATED MASS STORE, NEEDS 14 TO 17 ALSO)
C 14 RANMF     RANDOM NUMBER GENERATOR
C 15 NAIVE     DISCRETE FOURIER TRANSFORM COMPUTED NAIVELY
C 16 MFREAD   DUMMY I/O ROUTINE USING SIMULATED MASS STORE
C 17 MFWRIT   AS ABOVE
C
C 18 CYBER MASS STORE SAMPLE PROGRAM
C 19 MFREAD/MFWRIT CYBER MASS STORE I/O ROUTINE
C
C 20 CYBER EXTENDED CORE/LCM SAMPLE PROGRAM
C 21 MFREAD/MFWRIT CYBER EXTENDED CORE/LCM I/O ROUTINES
C
C 22 PDP 11 MASS STORE SAMPLE PROGRAM
C 23 MFREAD   PDP 11 STANDARD FORTRAN MASS STORE I/O ROUTINES
C 24 MFWRIT   AS ABOVE
C
C 25 PDP 11 FAST MACRO I/O SAMPLE PROGRAM (NEEDS MACRO OPEN SUBRTN)
C 26 MFREAD/MFWRIT PDP 11 FAST MACRO I/O ROUTINE FOR RSX11M/RT11.
C
C   END
C
C   SUBROUTINE RMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX, IPAK)
C REAL-TO-COMPLEX FFT (OR VICE-VERSA) OF MULTI-DIMENSND MASS STORE ARRAY
C (FRASER, ACM TOMS - 1978/79, AND J.ACM, V.23,N 2, APRIL 76, PP. 298-309)

```

```

C MASS STORE ARRAY IS EITHER REAL OR COMPLEX DATA (SEE NOTE BELOW)
C
C NOTE WELL THAT TYPE COMPLEX DATA MUST EXIST AS ALTERNATING
C REAL/IMAG/REAL/IMAG... ELEMENTS, BOTH IN MASS STORE AND IN FORTRAN
C WORKING ARRAY BUFA; IN THIS FFT, DIFFERENT SUBROUTINES WILL SET
C DIFFERENT TYPE (REAL OR COMPLEX) FOR ARRAY BUFA.
C
C MEXA(J) LIST OF DIMENSION SIZE EXPONS (BASE 2), ADJACENT VARIABLES FIRST
C NDIM IS NUMBER OF EXPONENTS IN LIST AND THUS THE NUMBER OF DIMENSIONS
C SUM TO NDIM OF MEXA(J) = M, WHERE 2**M IS EFFECT SIZE OF MASS STORE ARRAY
C   THUS, 2**M PACKED REAL VALUES,
C   OR, 2**M COMPLEX VALUES, IF COMPLEX RESULT WITH IPAK=1 (SEE BELOW)
C RMFFT ALWAYS REVERSES DIMENSION ORDER AND MEXA LIST (TRANSPOSED, IF 2 D)
C
C ISGN GIVES SIGN OF COMPLEX EXPONENT OF TRANSFORM (+ OR -), AND
C IDIR DETERMINES DIRECTION OF TRANSFORM, THUS:
C   IDIR=-1, REAL-TO-COMPLEX
C   IDIR=+1, COMPLEX-TO-REAL
C SCAL IS REAL MULTIPLIER OF RESULT (EG. SET SCAL=1. FWD, 1./2**M INV)
C
C BUFA IS CORE STORE WORKING ARRAY BASE ADDRESS (SEE NOTE ABOVE)
C IBEX, ICEX ARE BLOCK AND CORE SIZE EXPONENTS, THUS
C 2**IBEX IS NUMBER OF REAL ELEMENTS IN MASS STORE BLOCK
C 2**ICEX IS NUMBER OF REAL ELEMENTS IN CORE STORE BUFA
C
C IPAK IS ARRAY PACKING DETERMINATOR, THUS:
C IPAK=+1 EXPANDS COMPLEX ARRAY TO FULL REDUNDANCY (SAME AS SUBRTN CMFFT)
C IPAK=0 COMPUTES COMPLEX ARRAY OF JUST OVER HALF SIZE (COMMON METHOD)
C IPAK=-1 HOLDS COMPLEX ARRAY AT EXACTLY HALF SIZE (2**(M-1) CPLX ELMTS)
C
C MASS STORE ARRAY MUST BE OPEN FOR ACCESS BY USER SUBRTNS MFREAD/MFWRITE:
C EG. SUBRTN MFREAD(BUFA,NB,JB) AND MFWRITE(BUFA,NB,JB) TRANSFER ONE
C BLOCK, INDEX JB, BETWEEN MASS STORE AND CORE STORE BUFA, NB REALS
C (1.LE.JB.LE.2**(M-IBEX) IF REAL, OR 2**(M-IBEX+1) IF CPLX AND IPAK=1)
C
C   COMPLEX BUFA(1)
C   INTEGER B,MEXA(1)
C
C   MH=MFSUM(MEXA,NDIM,99)-1
C   B=IBEX-1
C   IF(IDIR.GT.0)GO TO 10
C
C
C BELOW, REAL-TO-COMPLEX TRANSFORM
C   MEXA(1)=MEXA(1)-1
C   CALL CMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX)
C   CALL MFRCMP (MEXA,NDIM,ISGN,IDIR,IPAK, BUFA,B,MH)
C   MEXA(NDIM)=MEXA(NDIM)+1
C   RETURN
C
C BELOW, COMPLEX-TO-REAL TRANSFORM
10  MEXA(NDIM)=MEXA(NDIM)-1
C   CALL MFRCMP (MEXA,NDIM,ISGN,IDIR,IPAK, BUFA,B,MH)
C   CALL CMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX)
C   MEXA(1)=MEXA(1)+1
C   RETURN
C
C   END

SUBROUTINE CMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX)
C
C COMPLEX FFT OF MULTI-DIMENSND MASS STORE ARRAY, CALLED BY USER OR RMFFT
C FOR COMMENTS, SEE SUBRTN RMFFT; ARGUMNTS HAVE SAME MEANING EXCEPT FOR
C IDIR=+1 OR -1, ARRAY ALWAYS COMPLEX, DIMENSION ORDER REVERSED
C WHILE IDIR=0, DIMENSION ORDER (AND MEXA LIST) ARE NOT REVERSED
C
C   COMPLEX BUFA(1)
C   INTEGER B,C,MEXA(1)
C   DATA LSET/1/,LREAD/1/,LWRIT/2/
C
C   M=MFSUM(MEXA,NDIM,99)
C   MREAL=M+1
C   B=IBEX-1
C   C=ICEX-1
C   NC=2**C
C   IPAS=0

```

```

      MPAS=M-C
      NPAS=C-B
C MOST EFFICIENT USE OF CORE STORE - TRIES TO DO C-B PASSES PER LOAD
      IDUM=MFINDX(LSET,B,M,M,NPAS)
C DUMMY CALL TO MFINDX TO SPECIFY VIRTUAL (B.'S'.M)**(C-B) PERMUTATION
C
C FIRST, PIECE-MEAL ATTACK ON FFT COMPUTATION FOLLOWS
10  IF(MPAS.LE.0) GO TO 40
      IF(MPAS.LT.NPAS)NPAS=MPAS
20  CALL MFLOAD(LREAD,BUFA,IBEX,ICEX,IFLG)
C LOAD CORE WORKING SPACE WITH 2**(C-B) BLOCKS ACCORDING TO MFINDX
      IF(IFLG.LT.0) GO TO 30
      CALL MFCOMP (MEXA,NDIM,ISGN, BUFA,B,C,M, NPAS,IPAS)
C DO MODIFIED IN-CORE FFT, REQUIRING NPAS PASSES STARTING WITH IPAS
      CALL MFLOAD(LWRIT,BUFA,IBEX,ICEX,IFLG)
C UNLOAD CORE AREA, WRITING BLOCKS BACK IN-PLACE TO MASS STORE
      GO TO 20
30  IPAS=IPAS+NPAS
      MPAS=MPAS-NPAS
      GO TO 10
C END OF FIRST PART
C
C SPECIFY BLOCKS TO BE READ IN NEXT PART IN TRUE ORDER (NO PERM)
40  IDUM=MFINDX(LSET,B,M,M,0)
C FINAL, CONCLUDING ATTACK ON FFT COMPUTATION FOLLOWS
50  CALL MFLOAD(LREAD,BUFA,IBEX,ICEX,IFLG)
      IF(IFLG.LT.0)GO TO 80
      CALL MFCOMP (MEXA,NDIM,ISGN, BUFA,C,C,M, C,IPAS)
C DO FINAL, C-PASS IN-CORE FFT OF EACH CORE-LOAD
      IF(SCAL.EQ.1.)GO TO 70
      DO 60 J=1,NC
60  BUFA(J)=BUFA(J)*SCAL
70  CALL MFLOAD(LWRIT,BUFA,IBEX,ICEX,IFLG)
      GO TO 50
C
C BELOW, SORT ARRAY (FULL BIT-REVERSAL AND DIMEN REVERSAL IF IDIR.NE.0)
80  IF(IDIR.EQ.0)GO TO 90
      CALL MFSORT(BUFA,IBEX,ICEX,1,MREAL,MREAL)
      M=MFSUM(MEXA,NDIM,-1)
C DO FULL BIT-REVERSAL OF M BITS (AND REVERSE MEXA LIST)
      RETURN
C
C BELOW, REVERSE BITS OF EACH DIMEN SEPARATELY (NO DIMEN REVERSAL)
90  IH=1
      DO 100 J=1,NDIM
          IG=IH
          IH=IH+MEXA(J)
100 CALL MFSORT(BUFA,IBEX,ICEX,IG,IH,MREAL)
      RETURN
C
      END

      SUBROUTINE MFCOMP (MEXA,NDIM,ISGN, BUFA,B,C,M, NPAS,IPAS)
C
C MODIFIED, IN-CORE FFT OF 2**C ELEMENTS, NPAS PASSES STARTING WITH IPAS
C MEXA, NDIM, ISGN ,BUFA AND M HAVE SAME MEANING AS IN RMFFT COMMENTS
C B,C EQUIVALENT TO IBEX,ICEX EXCEPT HERE REFER TO NUM COMPLEX ELMTS
C (2**C CMPLX ELMTS IN CORE STORE BUFA, IN BLOCKS OF 2**B CMPLX ELMTS)
C
C FFT W PHASE FACTOR COMPUTED RECURSIVELY EXCEPT ON BLOCK BOUNDS
C MULTIDIMEN FFT ACHIEVED BY REPEATING W SEQUENCES
C
      INTEGER B,C,SPAN,STEP,MEXA(1)
      COMPLEX TEMP,W,D,BUFA(1)
      DATA LINDX/0/,LREST/4/
C
      DATA PI/3.141592653589793/
      PIMOD=PI*2.0**(1-M)
      IF(ISGN.LT.0)PIMOD=-PIMOD
      NC=2**C
C
C BELOW, NPAS COMPUTATION PASSES WHILE DATA IN-CORE
      DO 50 JPAS=1,NPAS
          KPAS=IPAS+JPAS-1
C KPAS IS GLOBAL COMPUTING PASS NUMBER (JPAS-1 IS LOCAL)

```

```

      KDIFF=M-MFSUM(MEXA,NDIM,KPAS)
      KPEFF=KPAS+KDIFF
C KPEFF IS EFFECTIVE GLOBAL PASS FOR MULTIDIMEN. FFT W GENERATION
      D=CEXP(CMPLX(0.,PIMOD*2.0**KPEFF))
C D IS USED FOR RECURSIVE MODIFICATION OF W PHASE FACTOR
C
      ITEM=C-JPAS
      SPAN=2**ITEM
      STEP=2*SPAN
C SPAN SEPARATES VALUES IN FFT KERNEL, STEP TO NEXT PAIR, SAME W
C
      IF(B.LT.ITEM)ITEM=B
      NB=2**ITEM
      IF(ITEM.GT.KDIFF)ITEM=KDIFF
      NRPT=2**ITEM
C NRPT COUNTS REPETITION OF W FOR MULTIDIMEN. FFT
C
      IMOD=2**(M-B-KPAS-1)
      IF(IMOD.LE.1)GO TO 20
      IDUM=MFINDX(LREST,0,0,0,0)
      MEXP=KPEFF
      ITEM=KDIFF-B
      IF(ITEM.GT.0)GO TO 10
      MEXP=B+KPAS
      ITEM=0
10    NCLR=2**ITEM
      PIMOD2=PIMOD*2.0**MEXP
C IMOD AND NCLR ARE USED TO COMPUTE W WITHOUT EXCEEDING SMALL INTEGER
C
C BELOW, START OF ONE PASS THROUGH CORE, NOTING BLOCK BOUNDARIES
20    DO 50 I1=1,SPAN,NB
      W=(1.,0.)
      IF(IMOD.LE.1)GO TO 30
      INDWM=MOD(MFINDX(LINDX,0,0,0,0)-1,IMOD)/NCLR
      ANDWM=INDWM
      W=CEXP(CMPLX(0.,PIMOD2*ANDWM))
C NEW W COMPUTED DIRECTLY AT BEGINNING OF NEW BLOCK AREA
C
C BELOW, COMPUTATIONS WITHIN EACH BLOCK OF 2**B CMPLX ELMTS
30    DO 50 I2=1,NB,NRPT
C
C BELOW, REPETITION OF SAME W DUE TO MULTIDIMEN FFT
      DO 40 I3=1,NRPT
      I4=I1+I2+I3-2
C
C BELOW, STEPPING THOUGH INDICES HAVING SAME W IN ONE DIMEN FFT
      DO 40 J=I4,NC,STEP
      K=J+SPAN
      TEMP=(BUFA(J)-BUFA(K))*W
      BUFA(J)=BUFA(J)+BUFA(K)
40    BUFA(K)=TEMP
C FFT 2-POINT KERNEL ARITHMETIC (ALGORITHM BIT-REVERSAL FOLLOWS COMPUT)
C
      W=W*D
C RECURSIVE MODIFICATION OF W WITHIN BLOCK BOUNDARIES
50    CONTINUE
      RETURN
C
      END

      SUBROUTINE MFSORT(BUFA,IBEX,ICEX,IG,IH,M)
C
C BIT-REVERSED PERMUTATION (IG.'R'.IH) OF MASS STORE REAL ARRAY
C REVERSES IH-IG BITS IN INDEX (M-1,...,IH-1,...,IG,...,0)
C NOTE THAT THIS IS MORE GENERAL THAN THE FULL M-BIT REVERSAL
C OF REFERENCE (FRASER, J.ACM, V.23, N.2, APR. 76, P. 306),
C BUT ALGORITHM IS LOGICALLY THE SAME, WITH ALTERED BIT LIMITS.
C BUFA,IBEX,ICEX AND M HAVE SAME MEANING AS IN COMMENTS IN RMFFT
C (BLOCKS 2**IBEX, CORE BUFA 2**ICEX, TOTAL 2**M, ALL REAL)
C
      REAL BUFA(1)
      DATA LSET/1/,LPERM/2/,LREAD/1/,LWRIT/2/
C
      IDUM=MFINDX(LSET,IBEX,M,M,0)
C DUMMY CALL TO INITIALISE MFINDX (INITIALLY UNPERMUTED ARRAY)

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      IF(IH-IG.LE.1)RETURN
      IF(IG.GE.IBEX)GO TO 50
      IF(IH.LE.ICEX)GO TO 60
C CHECK FOR SPECIAL CASES, REQUIRING SIMPLER TREATMENT
C
C BELOW, MIXED PERMUTATION OF BOTH ELEMENTS AND BLOCKS
      IPAS=0
      NPAS=ICEX-IBEX
C MOST EFFICIENT USE OF CORE STORE - TRIES TO DO ICEX-IBEX PASSES PER LOAD
      MPAS=IH-IBEX
      IF(IBEX-IG.LT.MPAS)MPAS=IBEX-IG
C
C BELOW, FIRST VIRTUAL 'S' PERMUTATIONS
10  IF(MPAS.LE.0)GO TO 40
      IF(MPAS.LT.NPAS) NPAS=MPAS
      IGCOR=IG+IPAS
      IF((IGCOR.GT.IBEX-1).AND.(IGCOR.GT.IBEX+NPAS-1))GO TO 30
      IF((IPAS.EQ.0).AND.(IGCOR.GT.IBEX+NPAS-1))GO TO 30
C BYPASS UNNECESSARY CORE LOAD IF TRIVIAL CASES
      IDUM=MFINDX(LPERM,IBEX,IH,M,NPAS)
C DUMMY CALL TO MFINDX TO SPECIFY VIRTUAL (IBEX.S.IH)**NPAS PERM
C
C BELOW, LOAD CORE ACCORDING TO VIRTUAL PERMUTATION AND PERM ELMTS
20  CALL MFLOAD(LREAD,BUFA,IBEX,ICEX,IPLG)
      IF(IPLG.LT.0)GO TO 30
      IF(IPAS.NE.0) CALL MFREV(BUFA,IGCOR,IBEX,ICEX,-1)
      CALL MFREV(BUFA,IGCOR,IBEX+NPAS,ICEX,-1)
C CARRY OUT IN-CORE, SYMMETRIC R PERMS. (ONE ONLY ON FIRST PASS)
      CALL MFLOAD(LWRIT,BUFA,IBEX,ICEX,IPLG)
C UNLOAD CORE AREA, WRITING BLOCKS BACK IN-PLACE TO MASS STORE
      GO TO 20
C
30  IPAS=IPAS+NPAS
      MPAS=MPAS-NPAS
      GO TO 10
C END OF FIRST PART
C
C BELOW, FINAL 'R' PERM. OF BLOCKS IN MASS STORE, IF (IH-2*IBEX+IG).GT.1
40  CALL MFREV(BUFA,0,IH-2*IBEX+IG,M-IBEX,IBEX)
      RETURN
C
C BELOW, PERMUTATION OF BLOCKS ONLY REQUIRED
50  CALL MFREV(BUFA,IG-IBEX,IH-IBEX,M-IBEX,IBEX)
      RETURN
C
C BELOW, PERMUTATION OF ELEMENTS IN CORE ONLY REQUIRED
60  CALL MFLOAD(LREAD,BUFA,IBEX,ICEX,IPLG)
      IF(IPLG.LT.0)RETURN
      CALL MFREV(BUFA,IG,IH,ICEX,-1)
      CALL MFLOAD(LWRIT,BUFA,IBEX,ICEX,IPLG)
      GO TO 60
C
      END

      SUBROUTINE MFREV(BUFA,IG,IH,M,IBEX)
C
C BIT-REVERSED PERMUTATION OF RANDOMLY ADDRESSABLE ELEMENTS
C REVERSES IH-IG BITS IN INDEX (M-1,...,IH-1,...,IG,...,0)
C (GENERAL PERM IG.'R'.IH, FRASER, J.ACM, V.23, N.2, APR 1976, P. 300)
C IF IBEX.LT.0, SORTS 2**M REAL ELMTS IN CORE BUFA,
C IF IBEX.GE.0, SORTS BLOCKS IN MASS STORE
C (2**IBEX REAL ELMTS PER BLOCK AND 2**M BLOCKS IN SECOND CASE)
C
C THE ALGORITHM MAINTAINS A SET OF 'REVERSED' INTEGERS IN ARRAY IRA()
C OF INCREASING NUMBER OF BITS, UP TO 2 LESS THAN (IH-IG) BITS.
C INCREMENTING A REVERSED INTEGER THEN REQUIRES THE ALTERNATE
C ADDITION OF NRA() TO IRA(), OR REPLACEMENT BY THE NEXT LOWER
C INCREMENTED REVERSED INTEGER IN THE HIERARCHY, RECURSIVELY.
C THIS IN ITSELF IS FAST, AS RECURSION DEPTHS ARE ON AVERAGE SMALL.
C BUT, IN ADDITION, ONLY QUARTER LENGTH SERIES ARE GENERATED (-2 BITS)
C AND THE FULL LENGTH DERIVED BY SCALING BY 2 AND ADDING OFFSETS.
C IN THIS FINAL STAGE, ONLY VALID SWAP PAIRS ARE GENERATED (1 OR 3 EACH)
C
C WITHIN THE INNER LOOPS, GROUPS OF 2**IG ELMTS ARE MOVED TOGETHER
C WHILE THIS IS REPEATED OVER 2**(M-IH) PARTS OF THE ARRAY,

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C CORRESPONDING TO THE UNPERMUTED BITS M TO IH AND IG-1 TO  $\emptyset$ .
C
  REAL BUFA(1),TEMP
  INTEGER IRA(16),NRA(16),IFOFA(3),IROFA(3)
C
  IHG=IH-IG-3
  IF(IHG.LE.-2)RETURN
C NO PERMUTATION REQUIRED
C
  NB=2**IBEX
  NB1=NB+1
  NG=2**IG
  NGDB=NG*2
  NH=2**IH
  NHHF=NH/2
C NG IS MOVEMENT GROUP SIZE, NH IS PERMUTATION REPLICATION SIZE
  NM=2**M
  NPARS=NM-NH+1
  NREV=NH/4
C
  DO 1 $\emptyset$  J=1,IHG
  IRA(J)= $\emptyset$ 
  NREV=NREV/2
1 $\emptyset$  NRA(J)=NREV
C REVERSED INTEGER RECURSION SETS INITIALISED
C
  NREV=NH/4
  IFOFA(1)=NG-1
  IROFA(1)=NH/2-1
  IFOFA(2)=-1
  IROFA(2)=-1
  IFOFA(3)=NH/2+NG-1
  IROFA(3)=NH/2+NG-1
C THREE PAIRS OF OFFSETS TO CONVERT QUARTER TO FULL LENGTH SERIES
C
  IFOR= $\emptyset$ 
  IREV= $\emptyset$ 
C
C BELOW, GENERATE INDEX PAIRS AND SWAP (IREV IS 'TOP' OF IRA() SET)
2 $\emptyset$  NOF=3
  IF(IFOR.GE.IREV)NOF=1
C SELECTS ONCE-ONLY SWAP PAIRS (EITHER 1 OR 3 PAIRS)
  DO 4 $\emptyset$  JOF=1,NOF
  IFOF=IFOFA(JOF)
  IROF=IROFA(JOF)
  DO 4 $\emptyset$  I1=1,NG
C REPETITION OVER GROUP OR SUPER ELEMENT OF NG ACTUAL ELEMENTS
  IN2F=IFOR+IFOF+I1
  IN2R=IREV+IROF+I1
  DO 4 $\emptyset$  I2=1,NPARS,NH
C REPETITION OF SAME PERMUTATION OVER ARRAY PARTS
  IN3F=IN2F+I2
  IN3R=IN2R+I2
  IF(IBEX.GE. $\emptyset$ )GO TO 3 $\emptyset$ 
C
C BELOW, IN-CORE ELEMENT SORTING
  TEMP=BUFA(IN3R)
  BUFA(IN3R)=BUFA(IN3F)
  BUFA(IN3F)=TEMP
  GO TO 4 $\emptyset$ 
C
C BELOW, SORTING WHOLE BLOCKS IN MASS STORE
3 $\emptyset$  CALL MFREAD(BUFA,NB,IN3F)
  CALL MFREAD(BUFA(NB1),NB,IN3R)
  CALL MFWRIT(BUFA,NB,IN3R)
  CALL MFWRIT(BUFA(NB1),NB,IN3F)
C
4 $\emptyset$  CONTINUE
C END OF INNER, REPETITION LOOPS
C
  IFOR=IFOR+NGDB
C INCREMENT FORWARD QUARTER-LENGTH INTEGER (ALREADY SCALED BY NG*2)
  IF(IFOR.GE.NHHF)RETURN
C RETURN FORM SUBROUTINE
C
  IF(IREV.GE.NREV)GO TO 5 $\emptyset$ 

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C TEST FOR ALTERNATE METHODS OF REVERSE-INCREMMENTING (SIMPLE BELOW)
C NOTE THAT REVERSE QUARTER-LENGTH INTEGER IS ALREADY SCALED BY NG*2
C
      IREV=IREV+NREV
      GO TO 20

C
C ALTERNATE RECURSIVE ALTERATION TO QUARTER-LENGTH REVERSED SERIES
50  DO 60 J=1,IHG
      IF(IRA(J).LT.NRA(J))GO TO 70
60  CONTINUE
C
C BELOW, SIMPLE INCREMENT OF REVERSE INTEGER, LOWER IN HIERARCHY
70  IRA(J)=IRA(J)+NRA(J)
      IREV=IRA(J)
80  IF(J.EQ.1)GO TO 20
      J=J-1
      IRA(J)=IREV
      GO TO 80

C
      END

      SUBROUTINE MFLOAD(LOAD,BUFA,IBEX,ICEX,IPLG)
C
C LOADS, UNLOADS CORE STORE ARRAY BUFA, 2**ICEX REALS, 2**IBEX PER BLOCK
C RETURNS IPLG=+1 NORMALLY, IPLG=-1 WHEN FINISHED ONE PASS OF MASS STORE
C BLOCKS INDEXED ACCORDING TO VIRTUAL PERMUTATION FUNCTION MFINDX
C LOAD=1 (LREAD) READS BLOCKS FROM MASS STORE INTO CORE STORE BUFA
C LOAD=2 (LWRIT) WRITES BLOCKS BACK IN-PLACE TO MASS STORE
C
      REAL BUFA(1)
      DATA LINDX/0/,LHOLD/3/,LREST/4/
C
      NB=2**IBEX
      NCB=2**(ICEX-IBEX)
      IF(LOAD.EQ.2)GO TO 30
      IPLG=+1
      IDUM=MFINDX(LHOLD,0,0,0,0)
C HOLDS CURRENT MFINDX VALUE FOR ENTRY 2 AND SUBRTN MFCOMP
      DO 10 J=1,NCB
          K=(J-1)*NB
          JB=MFINDX(LINDX,0,0,0,0)
          IF(JB.LT.0)GO TO 20
          CALL MFREAD(BUFA(K+1),NB,JB)
C READS BLOCK WITH NEXT VIRTUAL MFINDX INDEX
10  CONTINUE
      RETURN
20  IPLG=-1
      RETURN
C
30  IDUM=MFINDX(LREST,0,0,0,0)
C RESETS MFINDX TO START OF IN-PLACE BLOCK
      DO 40 J=1,NCB
          K=(J-1)*NB
          JB=MFINDX(LINDX,0,0,0,0)
          CALL MFWRITE(BUFA(K+1),NB,JB)
C WRITES BLOCK WITH NEXT VIRTUAL MFINDX INDEX (REPEAT MFREAD SEQUENCE)
40  CONTINUE
      RETURN
C
      END

      FUNCTION MFINDX(LSPEC,B,H,M,N)
C
C VIRTUAL 'S' PERMUTATION (FRASER, J.ACM, V.23, N.2, APR. 76, P.303)
C CYCLIC SHIFTS H-B BITS IN INDEX (M-1,...,H-1,...,B,...,0)
C COMPUTES NEXT INDEX FOR SEQUENTIAL CORE LOAD, PERM (B.'S'.H)**N
C BLOCK SIZE EXPON B, MASS STORE EXPON M (0.LE.B.LE.H.LE.M)
C N IS EFFECTIVE NUMBER OF LEFT SHIFTS PER I/O PASS (-N RIGHT SHIFTS)
C
C NOTE VARIABLE NAMES AS FOLLOWS:
C IPERM IS 'P' OF ALGORITHM
C NPERM=N (ARGUMENT) IS 'N' OF ALGORITHM
C ISTEP IS 'Q**P' OF ALGORITHM

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C      JAY AND KAY ARE 'J' AND 'K' OF ALGORITHM
C NOTE UPPER BOUND H INSTEAD OF M, REQUIRING 2**(M-H) REPEATS
C
C LSPEC=0 (LINDX) RETURNS MFINDX FOR INDEX (B,H,M,N DUMMIES HERE)
C LSPEC=1 (LSET) SETS IPERM=0 (UNPERMED), ENTERS B,H,M,N PARAMS
C LSPEC=2 (LPERM) CHANGES THE B,H,M,N PARAMETERS
C LSPEC=3 (LHOLD) HOLDS CURRENT INDEXING STATE (B,H,M,N DUMMIES HERE)
C LSPEC=4 (LREST) RESTORES STATE TO LAST LHOLD (B,H,M,N DUMMIES HERE)
C
C      INTEGER B,H
C
C      IF(LSPEC.EQ.1)GO TO 100
C      IF(LSPEC.EQ.2)GO TO 200
C      IF(LSPEC.EQ.3)GO TO 300
C      IF(LSPEC.EQ.4)GO TO 400
C
C 100  IF(ISTEP.NE.0)GO TO 200
C BELOW, PRECEDES FIRST MFINDX OF A PASS
C      IF(NPERM.GT.0)IPERM=MOD(IPERM-NPERM,IHB)
C      IF(IPERM.LT.0)IPERM=IHB+IPERM
C      ISTEP=2**IPERM
C
C 200  BELOW, NORMAL GENERATION OF NEXT MFINDX
C      MFINDX=JAY+JOFF
C      IF(MFINDX.GT.NMB)GO TO 400
C      KAY=JAY+ISTEP
C      JAY=MOD(KAY,NHB)
C      IF(KAY.GE.NHB)JAY=JAY+1
C      NRPT=NRPT-1
C      IF(NRPT.GT.0)RETURN
C
C 300  NRPT,JOFF REQUIRED TO REPEAT SEQUENCE ON 2**(M-H) PARTS OF ARRAY
C      JOFF=JOFF+NHB
C      NRPT=NHB
C      JAY=0
C      RETURN
C
C 400  BELOW, END OF ONE PASS, PARS RESET, IPERM ALTERED IF INVERSE
C      IF(NPERM.LT.0)IPERM=MOD(IPERM-NPERM,IHB)
C      MFINDX=-1
C      JOFF=1
C      ISTEP=0
C      GO TO 300
C
C LSPEC=1 (LSET) SETS IPERM=0 (UNPERMED), ENTERS B,H,M,N PARAMS
100  IPERM=0
C
C LSPEC=2 (LPERM) CHANGES THE B,H,M,N PARAMETERS (DUMMIES ELSEWHERE)
200  IHB=H-B
C      NPERM=N
C      NHB=2**IHB
C      NMB=2**(M-B)
C      MFINDX=IPERM
C      GO TO 500
C
C LSPEC=3 (LHOLD) HOLDS CURRENT MFINDX INDEXING PARAMETERS
300  JAYH=JAY
C      JOFH=JOFF
C      NRPTH=NRPT
310  MFINDX=IPERM
C      RETURN
C
C LSPEC=4 (LREST) RESTORES PARAMETERS TO INDEX MFINDX AT LAST LHOLD
400  JAY=JAYH
C      JOFF=JOFH
C      NRPT=NRPTH
C      GO TO 310
C
C      END

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FUNCTION MFSUM(MEXA,NDIM,MLIM)

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C
C SCANS MEXA LIST IN REVERSE ORDER, RETURNING (MFSUM.JUST GT.MLIM)
C (IF MLIM LARGE ENOUGH, RETURNS M TOTAL FOR NDIM VALUES)

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C (IF MLIM NEGATIVE, RETURNS M TOTAL, REVERSES ORDER OF MEXA LIST)
C
C     INTEGER MEXA(1)
C
C     MFSUM=0
C     IF (NDIM.LE.0) RETURN
C
C     DO 10 J=1,NDIM
C     I=NDIM+1-J
C     MFSUM=MFSUM+MEXA(I)
C     IF ((MLIM.GE.0).AND.(MLIM.LT.MFSUM)) RETURN
10  CONTINUE
C     IF (MLIM.GE.0) RETURN
C
C BELOW, REVERSE ORDER OF MEXA LIST
C     NDIMH=NDIM/2
C     DO 20 J=1,NDIMH
C     K=NDIM+1-J
C     MTEM=MEXA(J)
C     MEXA(J)=MEXA(K)
20  MEXA(K)=MTEM
C     RETURN
C
C     END

SUBROUTINE MFRCMP (MEXA,NDIM,ISGN,IDIR,IPAK, BUFA,B,M)
C
C UNSCRAMBLES REAL-TO-COMPLEX FFT OR VICE-VERSA, CALLED BY SUBRTN RMFFT
C MOST ARGUMENTS HAVE SAME MEANING AS IN RMFFT COMMENTS
C BUT 2**B COMPLEX ELMTS IN MASS STORE BLOCK,
C USES (2**B)*4 CMLX IN BUFA, 'LOWER', 'UPPER' PLUS EXPANSION AREAS
C TOTAL MASS STORE ARRAY SIZE OF 2**M COMPLEX ELMTS.
C
C     COMPLEX ATEM,BTEM,TEMP,W,D,BUFA(1)
C     INTEGER B,JAYA(4),KAYA(4),JWKA(4),KWKA(4),MEXA(1)
C JAYA,KAYA,JWKA,KWKA ALLOW UP TO 4 DIMENSIONS - INCREASE IF REQUIRED
C     DATA PI/3.141592653589793/
C     DATA LOWER/1/,LUPPR/2/,LCLR/-1/
C
C     DO 10 IDIM=1,NDIM
C     JAYA(IDIM)=0
C     KAYA(IDIM)=0
10  CONTINUE
C MULTIDIMEN. CONJUGATE-SYMMETRIC INDICES ZEROED
C
C     IEXPND=1
C     NB=2**B
C     NDB=N*2
C     JBOF=2** (M-B)
C     MAX=M-B-MEXA (NDIM)
C     IF (IDIR*IPAK.LT.0) MAX=M-B
C     JBMAX=2**MAX
C     IF (MAX.LT.0) JBMAX=1
C     IF (IPAK.LT.0) JBMAX=0
C JBMAX IS MAXIMUM BLOCK INDEX REQUIRED (DEPENDS ON IPAK)
C
C     IWFG=0
C     W=(1.,0.)
C     D=CEXP (CMLX(0.,PI*2.0** (-MEXA (NDIM))))
C     IF (ISGN.LT.0) D=CONJG(D)
C W IS COMPLEX PHASE FACTOR, D IS RECURSIVE MODIFIER OF W
C
C     JAY=0
C     KAY=0
C     IDIM=NDIM
30  IF (IDIM.EQ.1) GO TO 40
C     NUMD=2**MEXA (IDIM)
C     JWKA (IDIM)=JAY
C     KWKA (IDIM)=KAY
C     JAY=JAY*NUMD+JAYA (IDIM)
C     KAY=KAY*NUMD+KAYA (IDIM)
C     IDIM=IDIM-1
C     GO TO 30
C CONJUGATE-SYMMETRIC BASE INDICES COMPUTED FROM MULTIDIMEN. SET
C

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40 MEX1=MEXA(1)
C 2**MEX1 IS NUMBER OF VALUES ADJACENT IN FIRST DIMENSION
  IFLG=-1
  IF(MEX1.LE.B)GO TO 140
C
C BELOW, FIRST DIMEN. GREATER THAN BLOCK SIZE, MULTIPLE BLOCKS
  NBPDI=2**(MEX1-B)
  NBLCNT=NBPDI
  KINC=NB
  KBINC=NBLCNT-1
  IF(JAY.EQ.KAY)NBLCNT=NBLCNT/2
  JB=JAY*NBPDI+1
  KB=KAY*NBPDI+1
C JB AND KB ARE BLOCK INDEX PAIRS CONTAINING CONJUGATE ELEMENTS
  J1=0
  K1=0
C
50 NCNT=NB+1
60 CALL MFRLOD(LOWER,IOF,BUFA,NB,JB,JBMAY,JBOF,IDIR,NCNT)
C LOWER BLOCK LOADED
  IF(JB.GT.JBMAY)IEXPND=-1
  J2=J1+IOF
  JB=JB+1
  J3=0
  K3=0
  NEWBLK=NCNT-NB
  IF(IFLG.GE.0)GO TO 80
C
70 CALL MFRLOD(LUPPR,IOF,BUFA,NB,KB,JBMAY,JBOF,IDIR,IFLG)
C UPPER BLOCK LOADED
  IFLG=IFLG+1
  K2=K1+IOF
  KB=KB+KBINC
C FIRST TIME, UPPER BLOCK STEPS HIGH, FOLLOWING STEPS SMALL NEGATIVE
  KBINC=-1
C
80 J=J2+J3
  K=K2+K3
C J AND K INDEX CONJUGATE-SYMMETRIC PAIRS IN CORE
  JJ=J+NBDB
  KK=K+NBDB
  IF(IDIR.GT.0)GO TO 180
C
C BELOW, UNSCRAMBLING FOR REAL-TO-COMPLEX FFT
  TEMP=(BUFA(J)+CONJG(BUFA(K)))*0.5
  BTEM=BUFA(K)-CONJG(BUFA(J))
  BTEM=(CMLX(AIMAG(BTEM),REAL(BTEM)))*0.5*W
  ATEM=TEMP+BTEM
  BTEM=TEMP-BTEM
  BUFA(J)=ATEM
  IF(IEXPND.GT.0)BUFA(JJ)=BTEM
  IF(IWFG.EQ.0)GO TO 150
  BUFA(K)=CONJG(BTEM)
  IF(IEXPND.GT.0)BUFA(KK)=CONJG(ATEM)
C
90 J3=J3+1
  K3=KINC-J3
C IN-CORE INDEX PAIRS STEPPED IN OPPOSING DIRECTIONS
  IF(IDIM.NE.NDIM)GO TO 95
  IWFG=1
  W=W*D
C RECURSIVE MODIFICATION OF W IF UNIDIMEN. TRANSFORM
95 NCNT=NCNT-1
  IF(NCNT.LE.0)GO TO 100
C ENTER RECURSION ROUTINE IF OPERATION COMPLETE IN CURRENT DIMEN
  IF(J3.EQ.1)GO TO 70
  IF(NCNT.GT.NEWBLK)GO TO 80
C END OF INNER LOOP (NOTE SPECIAL CASE WHEN J3=1 ABOVE)
C
C BELOW, MAY REQUIRE TO READ NEW BLOCKS
  NBLCNT=NBLCNT-1
  IF(NBLCNT.GT.0)GO TO 50
  IF(JAY.EQ.KAY)GO TO 60
C JAY.EQ.KAY NEEDS SYMMETRICAL MIDDLE, OTHERWISE CURRENT DIMEN COMPLT
C
C 100 BELOW, RECURSION TO COMPUTE MULTIDIMEN. CONJUGATE-SYMMETRY

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100 JAYA(IDIM)=0
    KAYA(IDIM)=0
    IDIM=IDIM+1
    IF(IDIM.GT.NDIM)GO TO 120
    NUMD=2**MEXA(IDIM)
    IF(NUMD.LE.1)GO TO 120
    IF(IDIM.NE.NDIM)GO TO 105
    IWFC=1
    W=W*D
C RECURSIVE MODIFICATION OF W IF MULTIDIMEN. FFT
105 IF(JAYA(IDIM).EQ.0)GO TO 110
    IF((JWKA(IDIM)*NUMD+JAYA(IDIM)).EQ.
X (KWKA(IDIM)*NUMD+KAYA(IDIM)))GO TO 100
    IF(KAYA(IDIM).EQ.1)GO TO 100
C
110 JAYA(IDIM)=JAYA(IDIM)+1
    KAYA(IDIM)=NUMD-JAYA(IDIM)
C RECURSIVE STEPPING OF MULTIDIMEN. CONJUGATE-SYMMETRIC INDEX PAIRS
GO TO 20
C
C BELOW, OPERATION COMPLETE, TIDY UP AND RETURN FROM SUBROUTINE
120 DO 130 IAREA=1,2
    CALL MFRLOD(IAREA,IOF,BUFA,NB,LCLR,JBMAX,JBOF,IDIR,IFLG)
C DUMMY CALL TO MFRLOD TO WRITE ANY UNWRITTEN BLOCKS TO MASS STORE
130 CONTINUE
    RETURN
C RETURN FROM SUBROUTINE
C
C
C 140 BELOW, FIRST DIMEN. LESS THAN BLOCK SIZE, INDEX PAIRS ALL IN-CORE
140 NUMD1=2**MEX1
    NBPD1=NB/NUMD1
    NCNT=NUMD1
    KINC=NCNT
    KBINC=0
    IF(JAY.EQ.KAY)NCNT=NCNT/2+1
    JB=JAY/NBPD1+1
    KB=KAY/NBPD1+1
C JB AND KB ARE BLOCK INDEX PAIRS CONTAINING CONJUGATE ELEMENTS
    J1=(JAY-(JB-1)*NBPD1)*NUMD1
    K1=(KAY-(KB-1)*NBPD1)*NUMD1
    GO TO 60
C
C 150 BELOW, UNSCRAMBLING WITH W0 (IWFC=0) MUST BE TREATED DIFFERENTLY
150 IF(IEXPND.LT.0)GO TO 160
C BELOW, ARRAY EXPANSION (EITHER IPAK=+1 OR IPAK=0 AND STILL REDUNDANT)
    BUFA(K)=CONJG(ATEM)
    BUFA(KK)=CONJG(BTEM)
    GO TO 90
C 160 BELOW, NO ARRAY EXPANSION (EITHER IPAK=-1 OR IPAK=0 NOT REDUNDANT)
160 IF(J.EQ.K)GO TO 170
    BUFA(K)=CONJG(BTEM)
    GO TO 90
C 170 BELOW, SPECIAL CASE IF IPAK=-1 AND ELEMENTS ARE SAME
170 BUFA(J)=CMPLX(REAL(ATEM),REAL(BTEM))
    GO TO 90
C
C 180 BELOW, SCRAMBLING FOR COMPLEX-TO-REAL FFT
180 IF(IWFC.EQ.0)GO TO 200
    BTEM=CONJG(BUFA(K))
190 ATEM=(BUFA(J)+BTEM)
    BTEM=(BUFA(J)-BTEM)*W
    BTEM=CMPLX(AIMAG(BTEM),REAL(BTEM))
    BUFA(J)=ATEM-CONJG(BTEM)
    BUFA(K)=CONJG(ATEM)+BTEM
    GO TO 90
C
C 200 BELOW, SCRAMBLING WITH W0 (IWFC=0) MUST BE TREATED DIFFERENTLY
200 IF(IEXPND.LT.0)GO TO 210
    BTEM=BUFA(JJ)
    GO TO 190
C
C 210 BELOW, NO REDUNDANCY (EITHER IPAK=-1 OR IPAK=0 OR 1 NOT REDUND)
210 IF(J.EQ.K)GO TO 220
    BTEM=CONJG(BUFA(K))

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      GO TO 190
C
C 220 BELOW, SPECIAL CASE IF IPAK=-1 AND ELEMENTS ARE SAME
220 BTEM=CMPLX(AIMAG(BUFA(J)),0.)
      BUFA(J)=CMPLX(REAL(BUFA(J)),0.)
      GO TO 190
C
      END

      SUBROUTINE MFRLOD(IAREA,I0F,BUFA,NB,JB,JBMAX,JBOF,IDIR,NCNT)
C
C LOADS, UNLOADS CORE STORE ARRAY BUFA, FOR REAL FFT UNSCRAMBLING ROUTINE
C BLOCK SIZE NB CMPLX, BLOCK NUMBER JB (JB=-1 DOES FINAL TIDY O/P)
C JBMAX IS MAX BLOCK INDX FOR EXPANSN, JBOF OFFSET TO EXPANDING BLOCKS
C IDIR=-1 DIRECTION REAL/CMPLX, +1 CMPLX/REAL
C IAREA=1 (LOWER) OR 2 (UPPER) OF TWO AREAS IN LOGICAL UNSCRAMBLING
C NOTE THAT BLOCK NORMALLY PHYSICALLY LOADED IN THESE AREAS,
C BUT, IF BLOCK ALREADY RESIDENT, MAY BE IN DIFFERENT AREA, SO
C I0F RETURNED AS ACTUAL OFFSET IN BUFA TO LOADED BLOCK.
C USES (2**B)*4 CMPLX IN BUFA, 'LOWER', 'UPPER' PLUS EXPANSION AREAS
C RETURNS I0F AS BUFFER OFFSET TO AREA (MAY NOT BE SAME, IF BLOCK RESIDENT)
C
C NCNT IS COUNT OF ELEMENTS TO BE ACCESSED IN THIS LOAD, TO ALLOW
C NOTE TO BE TAKEN OF ANY PARTLY FILLED BLOCKS DURING EXPANSION,
C PREVENTING THE READING OF 'NON-EXISTENT' BLOCKS.
C LISTS JBPFA(), NCPFA() OF SIZE NPARF HOLD THIS INFORMATION, DEFAULTS
C TO ALL-READ IF EXCEEDED, BUT INCREASE NPARF ETC, IF PROBLEM.
C
      COMPLEX BUFA(1)
      INTEGER JBAREA(2),NCAREA(2),JBPFA(5),NCPFA(5)
      DATA JBAREA(1)/-1/,JBAREA(2)/-1/,NPARF/5/
C
      IF(JBAREA(1).GE.0)GO TO 200
C JBAREA() HOLDS INDEX OF BLOCK LOADED IN AREA 1 OR 2 (FIRST TIME -1 BELOW)
C
      IEXIST=-1
      DO 10 I=1,NPARF
10    JBPFA(I)=-1
C PRE-CLEAR PARTLY FILLED BLOCK LIST (ONCE BLOCK FILLED, ALSO CLEARED)
C
20    NBDB=NB*2
      IF(JB.LT.0)GO TO 50
      IF(IAREA.EQ.2)GO TO 30
      NCLOW=NCNT
      IF(MOD(NCNT,2).NE.0)NCLOW=(NCLOW-1)*2
30    NCHLD=NCLOW
      IF(IAREA.EQ.2)NCHLD=NCLOW-1
      IF(NCNT.LT.0)NCHLD=1
C NCHLD IS THE NUMBER OF ELEMENTS TO BE ACCESSED IN CURRENT READ
C
      DO 40 I=1,2
      IF(JB.EQ.JBAREA(I))GO TO 140
40    CONTINUE
C TEST DONE TO SEE IF REQUIRED BLOCK ALREADY IN CORE (TRIVIAL IF SO)
C
C OTHERWISE BELOW, FIRST WRITE OUT RESIDENT BLOCK, THEN READ IN NEW
50    I0F=(IAREA-1)*NB+1
C I0F IS BASE OFFSET OF CORE AREA WHERE BLOCK IS TO BE FOUND
      I0FDB=I0F+NBDB
      IF(JBAREA(IAREA).LT.0)GO TO 90
      CALL MFWRIT(BUFA(I0F),NBDB,JBAREA(IAREA))
C WRITE OUT BLOCK BEFORE READING NEW BLOCK
      IF(IDIR.GT.0)GO TO 90
      IF(JBAREA(IAREA).GT.JBMAX)GO TO 90
      IF(NCAREA(IAREA).GE.NB)GO TO 80
C
C BELOW, IF LAST BLOCK ONLY PART-FILLED, INDEX, ELMTS ACCESSED NOTED
      DO 60 I=1,NPARF
      IF(JBPFA(I).LT.0)GO TO 70
60    CONTINUE
C NO ROOM IN LISTS, DEFAULTS TO ALL READ
      I=NPARF
      IEXIST=I
70    JBPFA(I)=JBAREA(IAREA)
      RDM1=DIMA(1)

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      IF (IRMF.GE.Ø) RDM1=RDM1*2.
      TDM1=FSIZ/RDM1
      NCPFA(I)=NCAREA(IAREA)
C
8Ø    CALL MFWRIT(BUFA(IOFDB),NBDB,JBOF+JBAREA(IAREA))
C SIMILARLY, WRITE OUT BLOCK PAIR IF EXPANDING
C
C BELOW, READ BLOCK NOTING BLOCK INDX (READ EXPANDED, IF PART FILLED)
9Ø    JBAREA(IAREA)=JB
      IF (JB.LT.Ø) GO TO 13Ø
      CALL MFREAD(BUFA(IOF),NBDB,JB)
C READ REQUIRED BLOCK AND NOTE ACCESS COUNT
      NCAREA(IAREA)=NCHLD
      IF (JB.GT.JBMAX) GO TO 13Ø
      IF (IDIR.GT.Ø) GO TO 12Ø
C
C BELOW, EXPANSION - DOES BLOCK EXIST TO READ
      DO 1ØØ I=1,NPARF
      IF (JB.EQ.JBPFA(I)) GO TO 11Ø
1ØØ   CONTINUE
      IF (IEXIST.LT.Ø) GO TO 13Ø
11Ø   JBPFA(I)=-1
      NCAREA(IAREA)=NCPFA(I)+NCHLD
C IF BLOCK TO BE READ WAS ONLY PART FILLED, THEN IT EXISTS TO READ
12Ø   CALL MFREAD(BUFA(IOFDB),NBDB,JBOF+JB)
C READ EXPANDED BLOCK IF REQUIRED
C
13Ø   RETURN
C RETURN FROM SUBROUTINE
C
C BELOW, TRIVIAL CASE - BLOCK ALREADY LOADED
14Ø   IOF=(I-1)*NB+1
C IOF IS BASE OFFSET OF CORE AREA WHERE BLOCK IS TO BE FOUND
      IF (I.NE.IAREA) GO TO 15Ø
      NCAREA(I)=NCAREA(I)+NCHLD
C INCREASE ACCESS COUNT IF CURRENT IAREA MATCHES ORIGINAL IAREA
15Ø   RETURN
C
      END

      FUNCTION MFPAR(IRMF,ICOMP)
C
C HELPER ROUTINE TO CROSS-COMPUTE MASS STORE FFT FILE PARAMETERS
C PARAMETERS ARE HELD AND COMPUTED IN 3 COMMON AREAS (SEE BELOW)
C MFPAR RETURNS Ø NORMALLY, -1 IF NOT ALL MFINT CORRECT, +1 IBEX ERROR
C
C COMMON/MFARG/ HOLDS ARGUMENTS AS USED IN FFT CALLS, AS FOLLOWS:
C VARIABLE NAMES HAVE SAME MEANING AS COMMENTS, SUBROUTINE RMFFT
C MEXA() HOLDS EXPONENTS FOR UP TO 4 DIMENSIONS (R/T ZEROS EXCESS)
C NDIM NUM DIMENS, IBEX,ICEX BLOCK AND CORE EXPONS, IPAK RMFFT PACKING
C ISGN,IDIR,SCAL ARE IGNORED HERE, BUT INCLUDED FOR COMPLETENESS
C
C COMMON/MFVAL/,/MFINT/ RETURN COMPUTED VALUES, USEFUL FOR FILE ACCESS
C NDMA(4),DIMA(4) HOLD DIMENSION SIZES CORRESPONDING TO MEXA()
C (EG. NDMA(1)=DIMA(1)=2.**MEXA(1), ETC. AND =1. BEYOND NDIM)
C
C NTD1,TDM1 IS CURRENT TOTAL NUM OF 'RECDS' OF SIZE NRDI,RDMI
C NRDI,RDMI IS NUM OF REALS IN CURRENT FIRST DIMENSION
C (USEFUL FOR ACCESSING DATA BY MFREAD/MFWRIT, NRDI,RDMI REALS,
C ASSUMING THAT MFREAD/MFWRIT CAN HANDLE 'RECDS' OF DIFFERENT SIZES)
C
C NFBK,FBLK IS MAXIMUM FILE SIZE OF 'RECDS' OF SIZE NRBK,RBLK
C NTBK,TBLK IS CURRENT TOTAL NUM OF 'RECDS' OF SIZE NRBK,RBLK
C NRBK,RBLK IS NUM OF REALS IN FFT WORKING BLOCK (2**IBEX REALS)
C (GIVES MAX AND CURRENT FILE SIZE AND ACCESS BY FFT ROUTINES,
C NFBK.GT.NTBK ONLY WITH PACKED REAL DATA WHEN EXPANDING, IPAK=Ø OR 1)
C
C NRCR,RCOR IS NUM OF REALS IN FFT WORKING CORE (2**ICEX REALS)
C NSZE=SIZE=2.**M, WHICH IS THE EFFECTIVE TOTAL SIZE OF TRANSFORM,
C WHERE M IS SUM TO NDIM OF MEXA() (SEE RMFFT COMMENTS)
C
C NOTE THAT ALL /MFARG/ ARE INTGS (EXCEPT SCAL), ALL /MFVAL/ REALS
C (/MFINT/ IS INTEGER CONVERSION OF /MFVAL/, ANY VALUE OF MFINT
C IS SET -1 IF TOO LARGE, BY FIXMAX, AND MFPAR RETURNED -1 AS FLAG,
C FIXMAX SET BY DATA STATEMENT TO 32767. HERE, BUT ALTER TO SUIT)
C

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C NOTE WELL THAT TYPE COMPLEX DATA MUST EXIST AS ALTERNATING
C REAL/IMAG/REAL/IMAG... ELEMENTS, BOTH IN MASS STORE AND IN FORTRAN
C WORKING ARRAY BUFA; IN THIS FFT, DIFFERENT SUBROUTINES WILL SET
C ROUTINE ARGUMENTS HAVE THE FOLLOWING EFFECT:
C IRMF=-1, DATA IS PACKED REAL, +1 DATA IS COMPLEX, ROUTINE RMFFT,
C IRMF=0, DATA IS COMPLEX, ROUTINE CMFFT
C
C ICOMP=0, COMPUTES VALUES IN /MFVAL/ FROM VALUES GIVEN IN /MFARG/
C ICOMP=1, REVERSE COMPUTES EXPONENTS IN /MFARG/ FROM /MFVAL/
C (DIMA(),RBLK,RCOR GIVEN INSTEAD OF MEXA(),IBEX,ICEX)
C
C NOTE, ROUTINE FORCES ICEX, IBEX TO CORRECT RANGE, MFPAR=+1 IF CANNOT
C
COMMON/MFARG/MEXA(4),NDIM,ISGN,IDIR,SCAL,IBEX,ICEX,IPAK
COMMON/MFVAL/DIMA(4),TDM1,RDM1,FBLK,TBLK,RBLK,RCOR,SIZE
COMMON/MFINT/NDMA(4),NTD1,NRD1,NFBK,NTBK,NRBK,NRCR,NSZE
REAL VAL(11)
INTEGER INT(11)
EQUIVALENCE (VAL(1),DIMA(1)),(INT(1),NDMA(1))
C
DATA FIXMAX/32767./,NMAX/4/
C
MFPAR=0
IF(ICOMP.EQ.0)GO TO 20
ALG2=ALOG(2.)
IBEX=IFIX(ALOG(RBLK)/ALG2+0.5)
ICEX=IFIX(ALOG(RCOR)/ALG2+0.5)
C
DO 10 I=1,NDIM
10 MEXA(I)=IFIX(ALOG(DIMA(I))/ALG2+0.5)
C
20 M=0
DO 30 I=1,NMAX
IF(I.GT.NDIM)MEXA(I)=0
M=M+MEXA(I)
30 DIMA(I)=2.**MEXA(I)
SIZE=2.**M
C
IF(IRMF.EQ.0)GO TO 90
IF(ICEX.GT.M)ICEX=M
IF(IBEX.GT.ICEX-2)IBEX=ICEX-2
IF(IBEX.LT.2)MFPAR=1
C FORCES ICEX.NGT.M AND IBEX.NGT.ICEX-2, OR MFPAR=1 (IRMF=+/- 1)
C
40 RBLK=2.**IBEX
RCOR=2.**ICEX
C
FADD=SIZE
IF(IRMF.EQ.0.OR.IPAK.GT.0)GO TO 50
C FADD IS ADDITIONAL FILE SIZE IN REALS, 'SIZE' IF CMFFT OR IPAK=1
C
FADD=0.
IF(IPAK.LT.0)GO TO 50
C IPAK=-1 REQUIRES NO FILE EXPANSION
C
IDIM=NDIM
IF(IRMF.LT.0)IDIM=1
FADD=SIZE*2./DIMA(IDIM)
C FADD COMPUTED FOR PARTICULAR CASE OF IPAK=0, WHEN COMPLEX
C
50 FSIZ=SIZE+FADD
ITEM=IFIX(FSIZ/RBLK+0.5)
IF(FLOAT(ITEM)*RBLK+0.5.LT.FSIZ)ITEM=ITEM+1
FBLK=FLOAT(ITEM)
C FBLK IS MAXIMUM NUMBER OF 'RECDs', SIZE RBLK, POSSIBLE
TBLK=FBLK
IF(IRMF.GE.0)GO TO 60
C GENERALLY TBLK=FBLK, BUT FOR PACKED REAL NOT SO, BELOW
FSIZ=SIZE
TBLK=FSIZ/RBLK
C
60 TDM1=1.
RDM1=FSIZ
IF(NDIM.EQ.1)GO TO 70
C JOB COMPLETED IF NDIM=1
C

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C OTHERWISE COMPUTE TDMI AS NUMBER OF 'RECDS', SIZE RDM1 REALS
C
70 DO 80 I=1,11
    INT(I)=-1
    IF (VAL(I).LE.FIXMAX) INT(I)=IFIX(VAL(I)+0.5)
    IF (INT(I).LT.0.AND.MFPAR.EQ.0) MFPAR=-1
80 CONTINUE
C CONVERT VALUES IN /MFVAL/ TO INTEGERS IN /MFINT/ (-1 IF TOO LARGE)
RETURN
C
90 IF (ICEX.GT.M+1) ICEX=M+1
    IF (IBEX.GT.ICEX-1) IBEX=ICEX-1
    IF (IBEX.LT.1) MFPAR=1
    GO TO 40
C FORCES ICEX.NGT.M+1 AND IBEX.NGT.ICEX-1, OR MFPAR=1 (IRMF=0)
C
    END

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SUBROUTINE DMPERM (MEXA,NDIM,NSHFT,IREX, BUFA,IBEX,ICEX)

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C
C SHIFTS ORDER OF DIMENSIONS OF REAL OR COMPLEX MASS STORE ARRAY
C NOTE, THIS IS NOT USED BY FFT SUBRTNS BUT IS INCLUDED FOR COMPLETENESS
C (FRASER, ACM TOMS - 1978/79, AND J.ACM, V.23,N.2, APRIL 76, PP. 298-309)
C
C MEXA(J) LIST OF DIMENSION SIZE EXPONS (BASE 2), ADJACENT VARIABLES FIRST
C NDIM IS NUMBER OF EXPONENTS IN LIST AND THUS THE NUMBER OF DIMENSIONS
C SUM TO NDIM: MEXA(J)=M, WHERE 2**M IS SIZE OF MASS STORE ARRAY (SEE BELOW)
C NSHFT IS DIMENSION SHIFT COUNT, THUS:
C NSHFT=0, NO SHIFT OR CHANGE OCCURS
C NSHFT=1,2 ETC., FIRST TO NEXT DIMENSION, CIRC NSHFT PLACE SHIFT (MOD NDIM)
C NSHFT=-1, REVERSES THE ORDER OF DIMENSIONS
C IREX=0 REAL, 1 COMPLEX (THAT IS, MOVEMENT GROUP IS 2**IREX REALS,
C AND TOTAL MASS STORE SIZE IS 2**(M+IREX) REAL ELEMENTS)
C
    REAL BUFA(1)
    INTEGER MEXA(1)
C
    NS=MOD(NSHFT,NDIM)
    IF (NS.EQ.0) RETURN
    M=MFSUM(MEXA,NDIM,-1)+IREX
C FINDS M TOTAL AND REVERSES MEXA LIST
    CALL MFSORT(BUFA,IBEX,ICEX,IREX,M,M)
C INITIAL OVERALL BIT-REVERSAL M BITS ABOVE IREX BITS
    IF (NSHFT.LT.0) GO TO 10
C
C BELOW, REVERSAL OF TWO PARTS, TO FORM REQUIRED SHIFT
    IH=MFSUM(MEXA,NS,-1)+IREX
    CALL MFSORT(BUFA,IBEX,ICEX,IREX,IH,M)
C REVERSE LOWER PART OF MEXA LIST AND LOWER PART OF ARRAY BITS
    CALL MFSORT(BUFA,IBEX,ICEX,IH,M,M)
C SEPARATELY REVERSE UPPER PART OF ARRAY BITS
    IH=MFSUM(MEXA(NSHFT+1),NDIM-NS,-1)
C REVERSE UPPER PART OF MEXA LIST
    RETURN
C RETURN FROM SUBROUTINE AFTER CYCLIC SHIFTS
C
C BELOW, SEPARATELY REVERSE OVER EACH DIMENSION (DIMEN REVERSAL)
10 IH=IREX
    DO 20 J=1,NDIM
        IC=IH
        IH=IH+MEXA(J)
20 CALL MFSORT(BUFA,IBEX,ICEX,IC,IH,M)
    RETURN
C
    END

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C
C THIS PROGRAM TESTS THE MASS STORE FFT BY COMPARISON WITH NAIVE DFT
C FFT PARAMETERS MAY BE ALTERED AT WILL (SEE COMMENTS)
C MASS STORE IS SIMULATED BY FORTRAN ARRAYS (SEE DUMMY I/O SUBROUTINES)
C PRINTING MAY BE COPIOUS, OR ONLY MAX DIFFERENCES (SEE COMMENTS)
C TEST OK IF MAX DIFFERENCES ARE NEAR ORDER OF MACHINE ROUND-OFF
C

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C DIFFERENT TYPE (REAL OR COMPLEX) FOR ARRAY BUFA.
C
COMMON/MFARG/MEXA(4),NDIM,ISGN,DIR,SCAL,IBEX,ICEX,IPAK
COMMON/MFVAL/DIMA(4),TDM1,RDM1,FBLK,TBLK,RBLK,RCOR,SIZE
COMMON/MFINT/NDMA(4),NTD1,NRD1,NFBK,NTBK,NRBK,NRCR,NSZE
C
C COMMON AREAS /MFARG/,/MFVAL/,/MFINT/ USEFUL FOR RUNNING MASS STORE FFT
C /MFARG/ HOLDS ARGUMENTS USED IN FFT CALLES, MOSTLY EXPONENTS
C HELPER ROUTINE MFPAR COMPUTES VALUES FROM /MFARG/ INTO
C /MFVAL/ (REALS AS SOME LARGE), /MFINT/ (INTEGER EQUIVALENTS IF POSS)
C OR CAN REVERSE-COMPUTE SOME /MFARG/ FROM /MFVAL/
C SEE COMMENTS, ROUTINE MFPAR.
C
COMMON/MASS/RMAS(1024)
COMPLEX ANAIV(512),BNAIV(512),CUBFA(512),CDIF
REAL RANDA(1024),BUFA(1024)
EQUIVALENCE (CUBFA(1),BUFA(1))
C
C COMMON/MASS/RMAS() USED BY ROUTINES MFREAD/MFWRITE TO SIMULATE MASS STORE
C ANAIV(), BNAIV() HOLD RESULT OF NAIVE DFT FOR COMPARISON
C BUFA(1)=CUBFA(1) IS WORKING AREA IN CORE STORE FOR FFT AND PROGRAM
C RANDA HOLDS PSEUDO RANDOM DATA USED IN TEST
C
LP=5
IPRINT=0
C LP IS PRINTER LOGICAL UNIT, IPRINT=+1 FOR COPIOUS PRINT,
C IPRINT=0 FOR MAX DIFFERENCES ONLY, -1 OVERALL MAX DIFFERENCE ONLY
C
M=5
C M SETS THE OVERALL ARRAY SIZE FOR AUTO IBEX,ICEX,MEXA,NDIM STEPPING
C FIXED VALUES CAN BE USED (SEE COMMENTS BELOW DO 100 STATEMENTS)
C
IRMF=-1
C IRMF=-1 REAL ROUTINE RMFFT TEST, 0 COMPLEX ROUTINE CMFFT TEST
C
ISGN=-1
IDIR=-1
IPAK=1
C FFT ARGS, ISGN=+/- 1, IDIR=-1 (RMFFT), IDIR=1 OR 0 (CMFFT)
C IPAK=1 OR 0 (RMFFT), -1 GIVES APPARENT FAILURES DUE TO SQUEEZED RESULT
C
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
IF(IPRINT.LE.0)WRITE(LP,910)IPRINT,IRMF
910 FORMAT(31H1MASS STORE FFT TEST - IPRINT =,I3,
X 36H (1=COPIOUS,0=MAX DIFFS,-1=OVERALL),,
X 9H IRMF =,I3,28H (0=CMFFT TEST,1=RMFFT TEST) /)
DIFMG=0.
C
C BELOW, DO 100 COMPUTES ALL POSSIBLE MEXA FOR 1,2 AND 3 DIMENSIONS
DO 100 NDIM=1,3
M2M=M-NDIM+1
IF(NDIM.EQ.1)M2M=1
DO 100 M2=1,M2M
M3M=M2M-M2+1
IF(NDIM.NE.3)M3M=1
DO 100 M3=1,M3M
IF(NDIM.EQ.1)MEXA(1)=M
IF(NDIM.EQ.2)MEXA(1)=M-M2
IF(NDIM.EQ.3)MEXA(1)=M-M2-M3
MEXA(2)=M2
MEXA(3)=M3
C REPLACE DO 100 SET BY FIXED MEXA LIST, IF DESIRED
C
IBEX=2
ICEX=4
C DUMMY IBEX, ICEX SO NO ERROR IN MFPAR BELOW
IERR=MFPAR(IRMF,0)
IF(IERR.NE.0)GO TO 1000
C CALL HELPER ROUTINE TO COMPUTE SIZES USED BY NAIVE DFT SUBRTN
C
M=MFSUM(MEXA,NDIM,99)
AR=RANMF(1)
C RESET RANDOM NUMBER GENERATOR AND COMPUTE M IN CASE NOT GIVEN
DO 10 J=1,NSZE
RANDA(J)=RANMF(-1)
10 ANAIV(J)=CMPLX(RANDA(J),0.)

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C LOAD RANDOM NUMBERS FOR FFT AND FOR NAIVE SUBROUTINE
C
      CALL NAIVE(ANAIV,BNAIV,NDMA(1),NDMA(2),NDMA(3),ISGN)
C NAIVE DFT SUBROUTINE CALLED TO COMPUTE 'SLOW' FOURIER TRANSFORM
C
C
      IF(IRMF.EQ.0)GO TO 20
C BELOW, SET LIMITS FOR STEPPING IBEX, ICEX, WHEN CALLING RMFFT
      IBEXL=2
      ICEXL=4
      ICEXM=M
      GO TO 30
C
C SETS DIFFERENT LIMITS FOR IBEX, ICEX FOR CMFFT BELOW (IRMF.EQ.0)
20  IBEXL=1
      ICEXL=2
      ICEXM=M+1
C
C
30  IF(ICEXL.GT.ICEXM)GO TO 1000
      DO 100 ICEX=ICEXL,ICEXM
      IBEXM=ICEX-2
      IF(IRMF.EQ.0)IBEXM=ICEX-1
      DO 100 IBEX=IBEXL,IBEXM
C IBEX AND ICEX COMPUTED; REPLACE DO 100 ABOVE BY FIXED IBEX,ICEX IF REQUD.
C
      IERR=MFPAR(IRMF,0)
      IF(IERR.NE.0)GO TO 1000
      NCNT=NRD1
C HELPER ROUTN, NTD1 TOTAL NUMB OF NRD1 REALS IN FIRST DIMENSION
C (NCNT IS NUMBER OF REAL ELMTS IN FIRST DIMENSION)
      SCAL=1.
      IF(IRMF.EQ.0)GO TO 50
C SWITCH FOR COMPLEX ROUTINE CMFFT AT 50, REAL RMFFT BELOW
C
      DO 40 JB=1,NTD1
      K=(JB-1)*NCNT
      DO 35 I=1,NCNT
      J=K+I
35  BUFA(I)=RANDA(J)
      CALL MFWRIT(BUFA,NRD1,JB)
40  CONTINUE
C LOAD RANDOM NUMBERS IN REAL ARRAY IN 'RECDs' OF FIRST DIMEN LENGTH
C (NOTE, THIS REQUIRES MFREAD/MFWRIT TO BE ABLE TO ACCEPT 'RECORDS'
C OF DIFFERENT LENGTH; OTHERWISE MUST USE NTBK AND NRBK HERE)
C
      CALL RMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX, IPAK)
C REAL MASS STORE ROUTINE RMFFT TO TRANSFORM ARRAY TO COMPLEX RESULT
GO TO 70
C
C BELOW, USE COMPLEX ROUTINE CMFFT, NCNT NUM OF CMLX ELMTS IN FIRST DIMEN
50  NCNT=NCNT/2
      DO 60 JB=1,NTD1
      K=(JB-1)*NCNT
      DO 55 I=1,NCNT
      J=K+I
55  CBUFA(I)=CMLX(RANDA(J),0.)
      CALL MFWRIT(BUFA,NRD1,JB)
60  CONTINUE
C LOAD REAL VALUES IN CMLX ARRAY IN 'RECDs' OF FIRST DIMEN LENGTH
C (NOTE, THIS REQUIRES MFREAD/MFWRIT TO BE ABLE TO ACCEPT 'RECORDS'
C OF DIFFERENT LENGTH; OTHERWISE MUST USE NTBK AND NRBK HERE)
C
      CALL CMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX)
C COMPLEX MASS STORE ROUTINE CMFFT TO TRANSFORM ARRAY TO COMPLEX RESULT
C
70  IF(IPRINT.LE.0)GO TO 75
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
      WRITE(LP,910)IPRINT,IRMF
      WRITE(LP,920)NDIM,ISGN,IDIR,IBEX,ICEX,IPAK,
      X (MEXA(J),J=1,NDIM)
920  FORMAT(8H NDIM =,I3,8H ISGN =,I3,8H IDIR =,I3,
      X 8H IBEX =,I3,8H ICEX =,I3,8H IPAK =,I3,
      X 10H MEXA() =,3I5)
      WRITE(LP,930)
930  FORMAT(8H INDEX,12X,3HFFT,24X,5HNAIVE,24X,4HDIFF)
C

```

```

75   IERR=MFPAR(-IRMF,0)
      IF(IERR.NE.0)GO TO 1000
      NCNT=NRD1/2
C HELPER ROUTN, NTD1 TOTAL NUMB OF NRD1 REALS IN FIRST DIMENSION
C (NOTE IRMF=+1 FOR DATA IN COMPLEX STATE, NCNT ELMTS)
C
C BELOW, COMPARES FFT COMPLEX RESULT WITH NAIVE RESULT
      DIFM=0.
      DO 80 JB=1,NTD1
        K=(JB-1)*NCNT
        CALL MFREAD(BUFA,NRD1,JB)
C READ 'RECDs' OF FIRST DIMEN LENGTH (RECOMPUTED ABOVE BY MFPAR)
C (NOTE, THIS REQUIRES MFREAD/MFWRITE TO BE ABLE TO ACCEPT 'RECORDS'
C OF DIFFERENT LENGTH; OTHERWISE MUST USE NTBK AND NRBK HERE)
C
      DO 80 I=1,NCNT
        J=K+I
        INDEX=J-1
        IF(IDIR.NE.0)CDIF=CBUFA(I)-BNAIV(J)
        IF(IDIR.EQ.0)CDIF=CBUFA(I)-ANAIV(J)
        DIF=ABS(REAL(CDIF))
        IF(DIF.GT.DIFM)DIFM=DIF
        DIF=ABS(AIMAG(CDIF))
        IF(DIF.GT.DIFM)DIFM=DIF
        IF(IPRINT.LE.0)GO TO 80
        IF(IDIR.NE.0)WRITE(LP,940)INDEX,CBUFA(I),BNAIV(J),CDIF
        IF(IDIR.EQ.0)WRITE(LP,940)INDEX,CBUFA(I),ANAIV(J),CDIF
940  FORMAT(LX,I5,3(2X,2E13.4))
80   CONTINUE
C BELOW, PRINT INTERMEDIATE DIFFERENCES
C
      IF(IPRINT.GE.0)WRITE(LP,950)DIFM,NDIM,ISGN,IDIR,IBEX,ICEX,IPAK,
X (MEXA(J),J=1,NDIM)
950  FORMAT(10H MAX DIFF ,E11.4,LX,
X 8H NDIM =,I3,8H ISGN =,I3,8H IDIR =,I3,
X 8H IBEX =,I3,8H ICEX =,I3,8H IPAK =,I3,
X 10H MEXA() =,3I5)
      IF(DIFM.GT.DIFMG)DIFMG=DIFM
C
C BELOW, INVERT ISGN AND IDIR FOR INVERSE TRANSFORM (COMPLEX-TO-REAL)
      ISGN=-ISGN
      IDIR=-IDIR
      SCAL=1./SIZE
      IF(IRMF.NE.0)
X CALL RMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX, IPAK)
C EITHER ROUTINE RMFFT INVERSE TRANSFORMS ARRAY TO PACKED REAL
C
      IF(IRMF.EQ.0)
X CALL CMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX)
C OR ROUTINE CMFFT INVERSE TRANSFORMS ARRAY
C
      IF(IPRINT.LE.0)GO TO 85
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
      WRITE(LP,910)IPRINT,IRMF
      WRITE(LP,920)NDIM,ISGN,IDIR,IBEX,ICEX,IPAK,
X (MEXA(J),J=1,NDIM)
      WRITE(LP,960)
960  FORMAT(8H INDEX,4X,8HFFT/2**M,6X,5HINPUT,10X,4HDIFF)
C
85   IERR=MFPAR(IRMF,0)
      IF(IERR.NE.0)GO TO 1000
      NCNT=NRD1
      IF(IRMF.EQ.0)NCNT=NCNT/2
C HELPER ROUTN, NTD1 TOTAL NUMB OF NRD1 REALS IN FIRST DIMENSION
C (NCNT IS NUMBER OF ELMTS IN FIRST DIMENSION, REAL OR CMLPX)
C
C BELOW, COMPARES FFT INVERSE RESULT WITH INITIAL RANDOM INPUT
      DIFM=0.
      DO 90 JB=1,NTD1
        K=(JB-1)*NCNT
        CALL MFREAD(BUFA,NRD1,JB)
C READ 'RECDs' OF FIRST DIMEN LENGTH (RECOMPUTED ABOVE BY MFPAR)
C (NOTE, THIS REQUIRES MFREAD/MFWRITE TO BE ABLE TO ACCEPT 'RECORDS'
C OF DIFFERENT LENGTH; OTHERWISE MUST USE NTBK AND NRBK HERE)
C
      DO 90 I=1,NCNT

```

```

      J=K+I
      INDEX=J-1
      IF(IRMF.NE.0)RM=BUFA(I)
      IF(IRMF.EQ.0)RM=REAL(CBUFA(I))
      DIF=ABS(RM-RANDA(J))
      IF(DIF.GT.DIFM)DIFM=DIF
      IF(IPRINT.GT.0)WRITE(LP,940)INDEX,RM,RANDA(J),DIF
90    CONTINUE
C PRINT INVERSE RESULTS (SHOULD BE SAME AS INITIAL RANDOM SET)
C
      IF(IPRINT.GE.0)WRITE(LP,950)DIFM,NDIM,ISGN,IDIR,IBEX,ICEX,IPAK,
X (MEXA(J),J=1,NDIM)
      IF(DIFM.GT.DIFMG)DIFMG=DIFM
      ISGN=-ISGN
      IDIR=-IDIR
C RESTORE ISGN AND IDIR FOR FORWARD TRANSFORM
100  CONTINUE
C
C BELOW, PRINT OVERALL MAXIMUM DIFFERENCE
      WRITE(LP,970)DIFMG,M,ISGN,IDIR,IPAK
970  FORMAT(18H OVERALL MAX DIFF ,E11.4,3X,
X 48H FOR ALL MEXA(1 TO 3 DIM),IBEX,ICEX,MEXA FOR M =,I3,
X 19H, ISGN,IDIR(+/-) =,2I4,8H IPAK =,I4)
      STOP
C
1000 WRITE(LP,980)IERR
980  FORMAT(25H IBEX FORCED TOO SMALL OR,
X 30H NOT ALL /MFINT/ CORRECT, IERR,I4)
      STOP
      END

```

```

      FUNCTION RANMF(J)
C
C RANDOM NUMBER GENERATOR FOR MASS STORE FFT TEST
C
      IF(J.GE.0)GO TO 20
C POSITIVE J CAUSES RESET OF INITIAL K
C NEGATIVE J MUST BE USED NORMALLY
C
      MODULO=2048
      FLMOD=2048.0
      DO 10 I=1,15
10    K=MOD(5*K,MODULO)
      Z=FLOAT(K)/FLMOD
      RANMF=Z
      RETURN
C
20    K=J
      RANMF=J
      RETURN
C
      END

```

```

      SUBROUTINE NAIVE(ANAIV,BNAIV,NJ,NK,NL,ISGN)
C
C NAIVE DISCRETE FOURIER TRANSFORM - 1 TO 3 DIMENSIONS
C USED TO TEST MASS STORE FFT, INPUT ARRAY ANAIV(NJ,NK,NL)
C RESULT RETURNED IN BOTH ARRAYS ANAIV AND BNAIV
C ANAIV DIMENSIONS IN INITIAL ORDER, BNAIV REVERSED ORDER
C NJ,NK,NL DIMENSIONING, ISGN SIGN OF COMPLEX EXPONENT OF FOURIER
C
      COMPLEX TEMP,ANAIV(NJ,NK,NL),BNAIV(NL,NK,NJ)
      DATA PI/3.141592653589793/
C
      PI2=PI*2.0
      IF(ISGN.LT.0)PI2=-PI2
C
      DO 20 JB=1,NJ
      AJB=FLOAT(JB-1)/FLOAT(NJ)
      DO 20 KB=1,NK
      AKB=FLOAT(KB-1)/FLOAT(NK)
      DO 20 LB=1,NL
      ALB=FLOAT(LB-1)/FLOAT(NL)
C

```

```

      TEMP=(0.,0.)
      DO 10 JA=1,NJ
      AJA=FLOAT(JA-1)*AJB
      DO 10 KA=1,NK
      AKA=FLOAT(KA-1)*AKB
      DO 10 LA=1,NL
      ALA=FLOAT(LA-1)*ALB
10    TEMP=TEMP+ANAIV(JA,KA,LA)*CEXP(CMPLX(0.,PI2*(AJA+AKA+ALA)))
C
20    BNAIV(LB,KB,JB)=TEMP
C
      DO 30 JA=1,NJ
      DO 30 KA=1,NK
      DO 30 LA=1,NL
30    ANAIV(JA,KA,LA)=BNAIV(LA,KA,JA)
      RETURN
C
      END

      SUBROUTINE MFREAD(BUFA,NB,JB)
C
C DUMMY SUBROUTINE TO SIMULATE RANDOM ACCESS MASS STORE READ
C READ BLOCK, INDEX JB, FROM MASS STORE TO BUFA, NB REAL VALUES
C COMMON ARRAY RMAS SIMULATES MASS STORE ARRAY
C
      COMMON/MASS/RMAS(1024)
      REAL BUFA(NB)
C
      IOF=(JB-1)*NB
      DO 10 I=1,NB
      K=IOF+I
10    BUFA(I)=RMAS(K)
      RETURN
C
      END

      SUBROUTINE MFWRITE(BUFA,NB,JB)
C
C DUMMY SUBROUTINE TO SIMULATE RANDOM ACCESS MASS STORE WRITE
C WRITE BLOCK, INDEX JB, FORM BUFA TO MASS STORE, NB REAL VALUES
C COMMON ARRAY RMAS SIMULATES MASS STORE ARRAY
C
      COMMON/MASS/RMAS(1024)
      REAL BUFA(NB)
C
      IOF=(JB-1)*NB
      DO 10 I=1,NB
      K=IOF+I
10    RMAS(K)=BUFA(I)
      RETURN
C
      END

      PROGRAM MASTOM(TAPE1,INPUT,OUTPUT,TAPE6=INPUT,TAPE5=OUTPUT)
C
C CONTROL DATA 6000 AND CYBER MASS STORE I/O FFT TEST PROGRAM
C
C NOTE WELL THAT TYPE COMPLEX DATA MUST EXIST AS ALTERNATING
C REAL/IMAG/REAL/IMAG... ELEMENTS, BOTH IN MASS STORE AND IN FORTRAN
C WORKING ARRAY BUFA; IN THIS FFT, DIFFERENT SUBROUTINES WILL SET
C DIFFERENT TYPE (REAL OR COMPLEX) FOR ARRAY BUFA.
C
C
C      COMMON/FFTCOM/LUN,MINDX(512)
C
C COMMON /FFTCOM/ HOLDS LOGICAL UNIT NUMBER FOR MASS STORE I/O
C ARRAY MINDX HOLDS RECORD INDICES FOR CYBER MASS STORE I/O
C
      COMMON/MFARG/MEXA(4),NDIM,ISGN,IDIR,SCAL,IBEX,ICEX,IPAK
      COMMON/MFVAL/DIMA(4),TDM1,RDM1,FBLK,TBLK,RBLK,RCOR,SIZE
      COMMON/MFINT/NDMA(4),NTD1,NRD1,NFBK,NTBK,NRBK,NRCR,NSZE
C

```

```

C COMMON AREAS /MFARG/,/MFVAL/,/MFINT/ USEFUL FOR RUNNING MASS STORE FFT
C /MFARG/ HOLDS ARGUMENTS USED IN FFT CALLES, MOSTLY EXPONENTS
C HELPER ROUTINE MFPAR COMPUTES VALUES FROM /MFARG/ INTO
C /MFVAL/ (REALS AS SOME LARGE), /MFINT/ (INTEGER EQUIVALENTS IF POSS)
C OR CAN REVERSE-COMPUTE SOME /MFARG/ FROM /MFVAL/
C SEE COMMENTS, ROUTINE MFPAR.
C
      COMPLEX CBUFA(4096)
      REAL BUFA(8192)
      EQUIVALENCE (CBUFA(1),BUFA(1))
C
C BUFA(1)=CBUFA(1) IS WORKING AREA IN CORE STORE FOR FFT AND PROGRAM
C
      LUN=1
C LUN IS LOGICAL UNIT FOR CYBER MASS STORE I/O
C
      LP=5
      IPRINT=0
C LP IS PRINTER LOGICAL UNIT, IPRINT=+1 COPIOUS, 0 MAX DIFFERENCES ONLY
C
      IRMF=-1
C IRMF=-1 REAL ROUTINE RMFFT TEST, 0 COMPLEX ROUTINE CMFFT TEST
C
      ISGN=-1
      IDIR=-1
      IPAK=1
C FFT ARGS, ISGN=+/- 1, IDIR=-1 (RMFFT), IDIR=1 OR 0 (CMFFT)
C IPAK=1, 0 OR -1 (RMFFT), NOT USED (CMFFT)
C
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
      IF(IPRINT.LE.0)WRITE(LP,910)IPRINT,IRMF
910 FORMAT(31H1MASS STORE FFT TEST - IPRINT =,I3,
X 25H (1=COPIOUS,0=MAX DIFFS),,
X 9H IRMF =,I3,28H (0=CMFFT TEST,1=RMFFT TEST) /)
      DIFMG=0.
C
      MEXA(1)=8
      MEXA(2)=6
      NDIM=2
      IBEX=9
      ICEX=13
C MORE FFT ARGS, 2**8 ROWS OF 2**6 ELMTS, 2 DIMEN,
C FFT MASS STORE BLOCKS 2**IBEX REALS, BUFA 2**ICEX REALS
C
      IERR=MFPAR(IRMF,0)
      IF(IERR.NE.0)GO TO 1000
      CALL OPENMS(LUN,MINDX,NFBK+1,0)
C CYBER 'READMS/WRITMS' MASS STORE OPENED WITH 'NUM REC'+1=NFBK+1
C (NOTE HELPER ROUTN MFPAR RETURNS NFBK AS MAXIMUM FILE SIZE)
C
      NCNT=NRBK
C HELPER ROUTINE, NCNT NUM OF ELMTS IN FFT BLOCK, NTBK TOTAL BLOCKS
C
      M=MFSUM(MEXA,NDIM,99)
      SCAL=1.
      VALU=0.
      VINC=1./SIZE
      IF(IRMF.EQ.0)GO TO 50
C SWITCH FOR COMPLEX ROUTINE CMFFT AT 50, REAL RMFFT BELOW
C
      DO 40 JB=1,NTBK
      DO 35 I=1,NCNT
      BUFA(I)=VALU
35 VALU=VALU+VINC
      CALL MFWRITE(BUFA,NRBK,JB)
40 CONTINUE
C LOAD RAMP FUNCTN IN REAL ARRAY IN 'RECDS' OF NRBK LENGTH
C (NOTE, CYBER MFREAD/MFWRITE (USING READMS/WRITMS) CANNOT ACCEPT 'RECDS'
C OF DIFFERENT LENGTH; OTHERWISE COULD USE NTD1 'RECDS' OF NRDI HERE)
C
      CALL RMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX, IPAK)
C REAL MASS STORE ROUTINE RMFFT TO TRANSFORM ARRAY TO COMPLEX RESULT
      GO TO 70
C
C BELOW, USE COMPLEX ROUTINE CMFFT, NCNT NUM OF CMPLX ELMTS IN FFT BLOCK
50 NCNT=NCNT/2

```

```

DO 60 JB=1,NTBK
DO 55 I=1,NCNT
  CBUFA(I)=CMPLX(VALU,0.)
55 VALU=VALU+VINC
  CALL MFWRIT(BUFA,NRBK,JB)
60 CONTINUE
C LOAD RAMP FUNCTN IN COMPLEX ARRAY IN 'RECDs' OF NRBK LENGTH
C (NOTE, CYBER MFREAD/MFWRITE (USING READMS/WRITEs) CANNOT ACCEPT 'RECDs'
C OF DIFFERENT LENGTH; OTHERWISE COULD USE NTD1 'RECDs' OF NRDI HERE)
C
  CALL CMFFT (MEXA,NDIM,ISGN,DIR,SCAL, BUFA,IBEX,ICEX)
C COMPLEX MASS STORE ROUTINE CMFFT TO TRANSFORM ARRAY TO COMPLEX RESULT
C
70 IF(IPRINT.LE.0)GO TO 75
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
  WRITE(LP,910)IPRINT,IRMF
  WRITE(LP,920)NDIM,ISGN,DIR,IBEX,ICEX,IPAK,
  X (MEXA(J),J=1,NDIM)
920 FORMAT(8H NDIM =,I3,8H ISGN =,I3,8H DIR =,I3,
  X 8H IBEX =,I3,8H ICEX =,I3,8H IPAK =,I3,
  X 10H MEXA() =,3I5)
  WRITE(LP,930)
930 FORMAT(8H INDEX,12X,3HFFT)
C
  IERR=MFPAR(-IRMF,0)
  IF(IERR.NE.0)GO TO 1000
  NCNT=NRBK/2
C HELPER ROUTN, NTBK TOTAL NUMB OF NRBK REALS IN FFT BLOCK
C (NOTE IRMF=+1 FOR DATA IN COMPLEX STATE, NCNT ELMTS)
C
C BELOW, PROGRAM CAN ACCESS FFT COMPLEX RESULT (AND FORM COMPARISON, IF KNOWN)
DO 80 JB=1,NTBK
  K=(JB-1)*NCNT
  CALL MFREAD(BUFA,NRBK,JB)
C READ 'RECDs' OF NRBK LENGTH, TOTALLING NTBK (RECOMPUTED BY MFPAR)
C (NOTE, CYBER MFREAD/MFWRITE (USING READMS/WRITEs) CANNOT ACCEPT 'RECDs'
C OF DIFFERENT LENGTH; OTHERWISE COULD USE NTD1 'RECDs' OF NRDI HERE)
C
DO 80 I=1,NCNT
  INDEX=J-1
  IF(IPRINT.LE.0)GO TO 80
  WRITE(LP,940)INDEX,CBUFA(I)
940 FORMAT(1X,I5,2X,2E13.4)
80 CONTINUE
C
C BELOW, INVERT ISGN AND DIR FOR INVERSE TRANSFORM (COMPLEX-TO-REAL)
75 ISGN=-ISGN
  DIR=-DIR
  SCAL=1./SIZE
  IF(IRMF.NE.0)
  X CALL RMFFT (MEXA,NDIM,ISGN,DIR,SCAL, BUFA,IBEX,ICEX, IPAK)
C EITHER ROUTINE RMFFT INVERSE TRANSFORMS ARRAY TO PACKED REAL
C
  IF(IRMF.EQ.0)
  X CALL CMFFT (MEXA,NDIM,ISGN,DIR,SCAL, BUFA,IBEX,ICEX)
C OR ROUTINE CMFFT INVERSE TRANSFORMS ARRAY
C
  IF(IPRINT.LE.0)GO TO 85
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
  WRITE(LP,910)IPRINT,IRMF
  WRITE(LP,920)NDIM,ISGN,DIR,IBEX,ICEX,IPAK,
  X (MEXA(J),J=1,NDIM)
  WRITE(LP,960)
960 FORMAT(8H INDEX,4X,8HFFT/2**M,6X,5HINPUT,10X,4HDIFF)
C
85 IERR=MFPAR(IRMF,0)
  IF(IERR.NE.0)GO TO 1000
  NCNT=NRBK
  IF(IRMF.EQ.0)NCNT=NCNT/2
C HELPER ROUTN, NTBK TOTAL NUMB OF NRBK REALS IN FFT BLOCK
C (NCNT IS NUMBER OF ELMTS IN FFT BLOCK, REAL OR CMPLX)
C
C BELOW, COMPARES FFT INVERSE RESULT WITH INITIAL RANDOM INPUT
  DIFM=0.
  VALU=0.
  DO 90 JB=1,NTBK

```



```

      CALL MFREAD(BUFA,NRBK,JB)
C READ 'RECDs' OF NRBK LENGTH, TOTALLING NTBK (RECOMPUTED BY MFPAR)
C (NOTE, CYBER MFREAD/MFWRIT (USING READMS/WRITMS) CANNOT ACCEPT 'RECDs'
C OF DIFFERENT LENGTH; OTHERWISE COULD USE NTD1 'RECDs' OF NRD1 HERE)
C
      DO 90 I=1,NCNT
      IF(IRMF.NE.0)RM=BUFA(I)
      IF(IRMF.EQ.0)RM=REAL(CBUFA(I))
      DIF=ABS(RM-VALU)
      IF(DIF.GT.DIFM)DIFM=DIF
      IF(IPRINT.GT.0)WRITE(LP,990)I,RM,VALU,DIF
990  FORMAT(1X,I5,3(2X,E13.4))
      VALU=VALU+VINC
90  CONTINUE
C PRINT INVERSE RESULTS (SHOULD BE SAME AS INITIAL RAMP)
C
      IF(IPRINT.GE.0)WRITE(LP,950)DIFM,NDIM,ISGN,IDIR,IBEX,ICEX,IPAK,
X (MEXA(J),J=1,NDIM)
950  FORMAT(10H MAX DIFF ,E11.4,1X,
X 8H NDIM =,I3,8H ISGN =,I3,8H IDIR =,I3,
X 8H IBEX =,I3,8H ICEX =,I3,8H IPAK =,I3,
X 10H MEXA() =,3I5)
      IF(DIFM.GT.DIFMG)DIFMG=DIFM
C
      STOP
C
1000 WRITE(LP,980)IERR
980  FORMAT(25H IBEX FORCED TOO SMALL OR,
X 30H NOT ALL /MFINT/ CORRECT, IERR,I4)
      STOP
      END

```

```

      SUBROUTINE MFREAD(BUFA,NB,JB)
C
C CONTROL DATA 6000 AND CYBER MASS STORE I/O ROUTINES FOR FFT
C
C (LOGICAL UNIT LUN IN COMMON/FFTCOM/ MUST HAVE BEEN OPENED
C PREVIOUSLY; EG. CALL OPENMS(LUN,INDXARRAY,NREC+1,0)
C WHERE NREC=2*(M-IBEX+1) IF IPAK=1 OR LESS IF IPAK=0 OR -1)
C
C SEE ALSO ALTERNATIVE SUBROUTINES USING EXTENDED CORE OR LCM
C (IN ADDITION, GETW, PUTW OR READM, WRITEM MACROS CAN BE USED)
C
C READ BLOCK, INDEX JB, FROM MASS STORE TO BUFA, NB REAL VALUES
C
      COMMON/FFTCOM/LUN,MINDX(512)
      REAL BUFA(NB)
C
      CALL READMS(LUN,BUFA,NB,JB)
      RETURN
C
      ENTRY MFWRIT
C WRITE BLOCK, INDEX JB, FROM BUFA TO MASS STORE, NB REAL VALUES
C
      CALL WRITMS(LUN,BUFA,NB,JB,-1,0)
      RETURN
C
      END

```

```

      PROGRAM LCMTOM(INPUT,OUTPUT,TAPE6=INPUT,TAPE5=OUTPUT)
C
C CONTROL DATA 6000 AND CYBER EXTENDED CORE/LCM FFT TEST PROGRAM
C
C NOTE WELL THAT TYPE COMPLEX DATA MUST EXIST AS ALTERNATING
C REAL/IMAG/REAL/IMAG... ELEMENTS, BOTH IN MASS STORE AND IN FORTRAN
C WORKING ARRAY BUFA; IN THIS FFT, DIFFERENT SUBROUTINES WILL SET
C DIFFERENT TYPE (REAL OR COMPLEX) FOR ARRAY BUFA.
C
C
      LEVEL 3, LBUFA
      COMMON/FFTCOM/LBUFA(32768)
C
C LBUFA IN EXTENDED CORE/LCM SIMULATES FAST MASS STORE

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C DIMENSION LBUFA AS LARGE AS NECESSARY FOR RESULTANT ARRAY
C
COMMON/MFARG/MEXA(4),NDIM,ISGN,IDIR,SCAL,IBEX,ICEX,IPAK
COMMON/MFVAL/DIMA(4),TDML,RDML,FBLK,TBLK,RBLK,RCOR,SIZE
COMMON/MFINT/NDMA(4),NTD1,NRD1,NFBK,NTBK,NRBK,NRCR,NSZE
C
C COMMON AREAS /MFARG/,/MFVAL/,/MFINT/ USEFUL FOR RUNNING MASS STORE FFT
C /MFARG/ HOLDS ARGUMENTS USED IN FFT CALLES, MOSTLY EXPONENTS
C HELPER ROUTINE MFPAR COMPUTES VALUES FROM /MFARG/ INTO
C /MFVAL/ (REALS AS SOME LARGE), /MFINT/ (INTEGER EQUIVALENTS IF POSS)
C OR CAN REVERSE-COMPUTE SOME /MFARG/ FROM /MFVAL/
C SEE COMMENTS, ROUTINE MFPAR.
C
COMPLEX CBUFA(4096)
REAL BUFA(8192)
EQUIVALENCE (CBUFA(1),BUFA(1))
C
C BUFA(1)=CBUFA(1) IS WORKING AREA IN CORE STORE FOR FFT AND PROGRAM
C
LP=5
IPRINT=0
C LP IS PRINTER LOGICAL UNIT, IPRINT=+1 COPIOUS, 0 MAX DIFFERENCES ONLY
C
IRMF=-1
C IRMF=-1 REAL ROUTINE RMFFT TEST, 0 COMPLEX ROUTINE CMFFT TEST
C
ISGN=-1
IDIR=-1
IPAK=1
C FFT ARGS, ISGN=+/- 1, IDIR=-1 (RMFFT), IDIR=1 OR 0 (CMFFT)
C IPAK=1, 0 OR -1 (RMFFT), NOT USED (CMFFT)
C
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
IF(IPRINT.LE.0)WRITE(LP,910)IPRINT,IRMF
910 FORMAT(31H1MASS STORE FFT TEST - IPRINT =,I3,
X 25H (1=COPIOUS,0=MAX DIFFS),,
X 9H IRMF =,I3,28H (0=CMFFT TEST,1=RMFFT TEST) /)
DIFMG=0.
C
MEXA(1)=8
MEXA(2)=6
NDIM=2
IBEX=7
ICEX=13
C MORE FFT ARGS, 2**8 ROWS OF 2**6 ELMTS, 2 DIMEN,
C FFT MASS STORE BLOCKS 2**IBEX REALS, BUFA 2**ICEX REALS
C
IERR=MFPAR(IRMF,0)
IF(IERR.NE.0)GO TO 1000
NCNT=NRD1
C HELPER ROUTINE, NCNT NUM OF ELMTS IN FIRST DIMEN, NTD1 TOTAL
C
M=MFSUM(MEXA,NDIM,99)
SCAL=1.
VALU=0.
VINC=1./SIZE
IF(IRMF.EQ.0)GO TO 50
C SWITCH FOR COMPLEX ROUTINE CMFFT AT 50, REAL RMFFT BELOW
C
DO 40 JB=1,NTD1
DO 35 I=1,NCNT
BUFA(I)=VALU
35 VALU=VALU+VINC
CALL MFWRIT(BUFA,NRD1,JB)
40 CONTINUE
C LOAD RAMP FUNCTN IN REAL ARRAY IN 'RECS' OF NRD1 LENGTH
C (NOTE, 'LCM' MFMREAD/MFWRIT CAN ACCEPT 'RECORDS' OF DIFFERENT LENGTH;
C SO CAN USE NTD1,NRD1 HERE INSTEAD OF NTBK,NRBK AS IN CYBER MASS STORE)
C
CALL RMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX, IPAK)
C REAL MASS STORE ROUTINE RMFFT TO TRANSFORM ARRAY TO COMPLEX RESULT
GO TO 70
C
C BELOW, USE COMPLEX ROUTINE CMFFT, NCNT NUM OF CMLX ELMTS IN FFT BLOCK
50 NCNT=NCNT/2
DO 60 JB=1,NTD1

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        DO 55 I=1,NCNT
        CBUFA(I)=CMLPX(VALU,0.)
55      VALU=VALU+VINC
        CALL MFWRIT(BUFA,NRD1,JB)
60      CONTINUE
C LOAD RAMP FUNCTN IN COMPLEX ARRAY IN 'RECDs' OF NRD1 LENGTH
C (NOTE, 'LCM' MFREAD/MFWRIT CAN ACCEPT 'RECORDS' OF DIFFERENT LENGTH;
C SO CAN USE NTD1,NRD1 HERE INSTEAD OF NTBK,NRBK AS IN CYBER MASS STORE)
C
        CALL CMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX)
C COMPLEX MASS STORE ROUTINE CMFFT TO TRANSFORM ARRAY TO COMPLEX RESULT
C
70      IF(IPRINT.LE.0)GO TO 75
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
        WRITE(LP,910)IPRINT,IRMF
        WRITE(LP,920)NDIM,ISGN,IDIR,IBEX,ICEX,IPAK,
X (MEXA(J),J=1,NDIM)
920    FORMAT(8H NDIM =,I3,8H ISGN =,I3,8H IDIR =,I3,
X 8H IBEX =,I3,8H ICEX =,I3,8H IPAK =,I3,
X 10H MEXA() =,3I5)
        WRITE(LP,930)
930    FORMAT(8H INDEX,12X,3HFFT)
C
        IERR=MFPAR(-IRMF,0)
        IF(IERR.NE.0)GO TO 1000
        NCNT=NRD1/2
C HELPER ROUTINE, NCNT NUM OF ELMTS IN FIRST DIMEN, NTD1 TOTAL
C (NOTE IRMF=+1 FOR DATA IN COMPLEX STATE, NCNT ELMTS)
C
C BELOW, PROGRAM CAN ACCESS FFT COMPLEX RESULT (AND FORM COMPARISON, IF KNOWN)
        DO 80 JB=1,NTD1
        K=(JB-1)*NCNT
        CALL MFREAD(BUFA,NRD1,JB)
C READ 'RECDs' OF NRD1 LENGTH, TOTALLING NTD1 (RECOMPUTED BY MFPAR)
C (NOTE, 'LCM' MFREAD/MFWRIT CAN ACCEPT 'RECORDS' OF DIFFERENT LENGTH;
C SO CAN USE NTD1,NRD1 HERE INSTEAD OF NTBK,NRBK AS IN CYBER MASS STORE)
C
        DO 80 I=1,NCNT
        INDEX=J-1
        IF(IPRINT.LE.0)GO TO 80
        WRITE(LP,940)INDEX,CBUFA(I)
940    FORMAT(1X,I5,2X,2E13.4)
80      CONTINUE
C
C BELOW, INVERT ISGN AND IDIR FOR INVERSE TRANSFORM (COMPLEX-TO-REAL)
75      ISGN=-ISGN
        IDIR=-IDIR
        SCAL=1./SIZE
        IF(IRMF.NE.0)
X CALL RMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX, IPAK)
C EITHER ROUTINE RMFFT INVERSE TRANSFORMS ARRAY TO PACKED REAL
C
        IF(IRMF.EQ.0)
X CALL CMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX)
C OR ROUTINE CMFFT INVERSE TRANSFORMS ARRAY
C
        IF(IPRINT.LE.0)GO TO 85
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
        WRITE(LP,910)IPRINT,IRMF
        WRITE(LP,920)NDIM,ISGN,IDIR,IBEX,ICEX,IPAK,
X (MEXA(J),J=1,NDIM)
        WRITE(LP,960)
960    FORMAT(8H INDEX,4X,8HFFT/2**M,6X,5HINPUT,10X,4HDIFF)
C
85      IERR=MFPAR(IRMF,0)
        IF(IERR.NE.0)GO TO 1000
        NCNT=NRD1
        IF(IRMF.EQ.0)NCNT=NCNT/2
C HELPER ROUTINE, NCNT NUM OF ELMTS IN FIRST DIMEN, NTD1 TOTAL
C (NCNT IS NUMBER OF ELMTS IN FFT BLOCK, REAL OR CMLPX)
C
C BELOW, COMPARES FFT INVERSE RESULT WITH INITIAL RANDOM INPUT
        DIFM=0.
        VALU=0.
        DO 90 JB=1,NTD1
        CALL MFREAD(BUFA,NRD1,JB)

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C READ 'RECDs' OF NRD1 LENGTH, TOTALLING NTD1 (RECOMPUTED BY MFPAR)
C (NOTE, 'LCM' MFREAD/MFWRITE CAN ACCEPT 'RECORDS' OF DIFFERENT LENGTH;
C SO CAN USE NTD1, NRD1 HERE INSTEAD OF NTBK, NRBK AS IN CYBER MASS STORE)
C
  DO 90 I=1, NCNT
    IF (IRMF.NE.0) RM=BUFA(I)
    IF (IRMF.EQ.0) RM=REAL(CBUFA(I))
    DIF=ABS(RM-VALU)
    IF (DIF.GT.DIFM) DIFM=DIF
    IF (IPRINT.GT.0) WRITE(LP, 990) I, RM, VALU, DIF
990  FORMAT(1X, I5, 3(2X, E13.4))
    VALU=VALU+VINC
90  CONTINUE
C PRINT INVERSE RESULTS (SHOULD BE SAME AS INITIAL RAMP)
C
  IF (IPRINT.GE.0) WRITE(LP, 950) DIFM, NDIM, ISGN, IDIR, IBEX, ICEX, IPAK,
X (MEXA(J), J=1, NDIM)
950  FORMAT(10H MAX DIFF ,E11.4, 1X,
X 8H NDIM =, I3, 8H ISGN =, I3, 8H IDIR =, I3,
X 8H IBEX =, I3, 8H ICEX =, I3, 8H IPAK =, I3,
X 10H MEXA() =, 3I5)
  IF (DIFM.GT.DIFMG) DIFMG=DIFM
C
  STOP
C
1000 WRITE(LP, 980) IERR
980  FORMAT(25H IBEX FORCED TOO SMALL OR,
X 30H NOT ALL /MFINIT/ CORRECT, IERR, I4)
  STOP
  END

  SUBROUTINE MFREAD(BUFA, NB, JB)
C
C CONTROL DATA 6000 AND CYBER EXTENDED CORE STORE I/O ROUTINES FOR FFT
C
C (EXTENDED CORE OR LCM TAKES PLACE OF MASS STORE -
C DIMENSION LBUFA AS LARGE AS NECESSARY FOR LARGE ARRAYS)
C
C ALTERNATIVELY, EXTENDED CORE USE COULD BE COMBINED WITH
C READM/WRITE MACROS TO ACCESS LARGER FILE WHEN NECESSARY
C (USE VIRTUAL MEMORY ALGORITHM - MANY SECTORS HELD IN
C EXTENDED CORE AND ONLY ACCESSED FROM DISC IF NOT PRESENT).
C
C READ BLOCK, INDEX JB, FROM EXTENDED CORE TO BUFA, NB REAL VALUES
C
  LEVEL 3, LBUFA
  COMMON/FFTCOM/LBUFA(32768)
  REAL BUFA(NB)
C
  CALL MOVLEV(LBUFA((JB-1)*NB+1), BUFA, NB)
  RETURN
C
  ENTRY MFWRITE
C WRITE BLOCK, INDEX JB, FROM BUFA TO EXTENDED CORE, NB REAL VALUES
C
  CALL MOVLEV(BUFA, LBUFA((JB-1)*NB+1), NB)
  RETURN
C
  END

C
C PDP 11 MASS STORE I/O FFT SAMPLE TEST PROGRAM
C
C NOTE WELL THAT TYPE COMPLEX DATA MUST EXIST AS ALTERNATING
C REAL/IMAG/REAL/IMAG... ELEMENTS, BOTH IN MASS STORE AND IN FORTRAN
C WORKING ARRAY BUFA; IN THIS FFT, DIFFERENT SUBROUTINES WILL SET
C DIFFERENT TYPE (REAL OR COMPLEX) FOR ARRAY BUFA.
C
C
  COMMON/FFTCOM/LUN
C
C COMMON /FFTCOM/ HOLDS LOGICAL UNIT NUMBER FOR MASS STORE I/O
C

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COMMON/MFARG/MEXA(4),NDIM,ISGN,IDIR,SCAL,IBEX,ICEX,IPAK
COMMON/MFVAL/DIMA(4),TDML,RDML,FBLK,TBLK,RBLK,RCOR,SIZE
COMMON/MFINT/NDMA(4),NTD1,NRD1,NFBK,NTBK,NRBK,NRCR,NSZE
C
C COMMON AREAS /MFARG/,/MFVAL/,/MFINT/ USEFUL FOR RUNNING MASS STORE FFT
C /MFARG/ HOLDS ARGUMENTS USED IN FFT CALLES, MOSTLY EXPONENTS
C HELPER ROUTINE MFPAR COMPUTES VALUES FROM /MFARG/ INTO
C /MFVAL/ (REALS AS SOME LARGE), /MFINT/ (INTEGER EQUIVALENTS IF POSS)
C OR CAN REVERSE-COMPUTE SOME /MFARG/ FROM /MFVAL/
C SEE COMMENTS, ROUTINE MFPAR.
C
COMPLEX CBUFA(1024)
REAL BUFA(2048)
EQUIVALENCE (CBUFA(1),BUFA(1))
C
C BUFA(1)=CBUFA(1) IS WORKING AREA IN CORE STORE FOR FFT AND PROGRAM
C
LUN=1
C LUN IS LOGICAL UNIT FOR PDP 11 MASS STORE I/O
C
LP=5
IPRINT=0
C LP IS PRINTER LOGICAL UNIT, IPRINT=+1 COPIOUS, 0 MAX DIFFERENCES ONLY
C
IRMF=-1
C IRMF=-1 REAL ROUTINE RMFFT TEST, 0 COMPLEX ROUTINE CMFFT TEST
C
ISGN=-1
IDIR=-1
IPAK=1
C FFT ARGS, ISGN=+/- 1, IDIR=-1 (RMFFT), IDIR=1 OR 0 (CMFFT)
C IPAK=1, 0 OR -1 (RMFFT), NOT USED (CMFFT)
C
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
IF(IPRINT.LE.0)WRITE(LP,910)IPRINT,IRMF
910 FORMAT(31H1MASS STORE FFT TEST - IPRINT =,I3,
X 25H (1=COPIOUS,0=MAX DIFFS),,
X 9H IRMF =,I3,28H (0=CMFFT TEST,1=RMFFT TEST) /)
DIFMG=0.
C
MEXA(1)=8
MEXA(2)=6
NDIM=2
IBEX=7
ICEX=11
C MORE FFT ARGS, 2**8 ROWS OF 2**6 ELMTS, 2 DIMEN,
C FFT MASS STORE BLOCKS 2**IBEX REALS, BUFA 2**ICEX REALS
C
IERR=MFPAR(IRMF,0)
IF(IERR.NE.0)GO TO 1000
CALL ASSIGN(LUN,'FFTEST.DAT',0)
NWD=NRBK*2
DEFINE FILE LUN(NFBK,NWD,U,INDX)
C PDP 11 MASS STORE OPENED WITH NUM REC=NFBK, NRBK*2 WORDS PER REC
C (NOTE HELPER ROUTN MFPAR RETURNS NFBK AS MAXIMUM FILE SIZE)
C
NCNT=NRBK
C HELPER ROUTINE, NCNT NUM OF ELMTS IN FFT BLOCK, NTBK TOTAL BLOCKS
C
M=MFSUM(MEXA,NDIM,99)
SCAL=1.
VALU=0.
VINC=1./SIZE
IF(IRMF.EQ.0)GO TO 50
C SWITCH FOR COMPLEX ROUTINE CMFFT AT 50, REAL RMFFT BELOW
C
DO 40 JB=1,NTBK
DO 35 I=1,NCNT
BUFA(I)=VALU
35 VALU=VALU+VINC
CALL MFWRIT(BUFA,NRBK,JB)
40 CONTINUE
C LOAD RAMP FUNCTN IN REAL ARRAY IN 'RECDS' OF NRBK LENGTH
C (NOTE, PDP 11 MFREAD/MFWRITE (USING FORTRAN I/O) CANNOT ACCEPT 'RECDS'
C OF DIFFERENT LENGTH; OTHERWISE COULD USE NTD1 'RECDS' OF NRD1 HERE)
C

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      CALL RMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX, IPAK)
C REAL MASS STORE ROUTINE RMFFT TO TRANSFORM ARRAY TO COMPLEX RESULT
      GO TO 70
C
C BELOW, USE COMPLEX ROUTINE CMFFT,  NCNT NUM OF CMPLX ELMTS IN FFT BLOCK
50  NCNT=NCNT/2
      DO 60 JB=1,NTBK
          DO 55 I=1,NCNT
              CBUFA(I)=CMPLX(VALU,0.)
55  VALU=VALU+VINC
          CALL MFWRITE(BUFA,NRBK,JB)
60  CONTINUE
C LOAD RAMP FUNCTN IN COMPLEX ARRAY IN 'RECDS' OF NRBK LENGTH
C (NOTE, PDP 11 MFREAD/MFWRITE (USING FORTRAN I/O) CANNOT ACCEPT 'RECDS'
C OF DIFFERENT LENGTH; OTHERWISE COULD USE NTD1 'RECDS' OF NRD1 HERE)
C
      CALL CMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX)
C COMPLEX MASS STORE ROUTINE CMFFT TO TRANSFORM ARRAY TO COMPLEX RESULT
C
70  IF(IPRINT.LE.0)GO TO 75
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
      WRITE(LP,910)IPRINT,IRMF
      WRITE(LP,920)NDIM,ISGN,IDIR,IBEX,ICEX,IPAK,
      X (MEXA(J),J=1,NDIM)
920  FORMAT(8H NDIM =,I3,8H ISGN =,I3,8H IDIR =,I3,
      X 8H IBEX =,I3,8H ICEX =,I3,8H IPAK =,I3,
      X 10H MEXA() =,3I5)
      WRITE(LP,930)
930  FORMAT(8H INDEX,12X,3HFFT)
C
      IERR=MFPAR(-IRMF,0)
      IF(IERR.NE.0)GO TO 1000
      IF(IERR.NE.0)GO TO 1000
      NCNT=NRBK/2
C HELPER ROUTN,  NTBK TOTAL NUMB OF NRBK REALS IN FFT BLOCK
C (NOTE IRMF=+1 FOR DATA IN COMPLEX STATE,  NCNT ELMTS)
C
C BELOW, PROGRAM CAN ACCESS FFT COMPLEX RESULT (AND FORM COMPARISON,  IF KNOWN)
      DO 80 JB=1,NTBK
          K=(JB-1)*NCNT
          CALL MFREAD(BUFA,NRBK,JB)
C READ 'RECDS' OF NRBK LENGTH,  TALLING NTBK (RECOMPUTED BY MFPAR)
C (NOTE, PDP 11 MFREAD/MFWRITE (USING FORTRAN I/O) CANNOT ACCEPT 'RECDS'
C OF DIFFERENT LENGTH; OTHERWISE COULD USE NTD1 'RECDS' OF NRD1 HERE)
C
      DO 80 I=1,NCNT
          INDEX=J-1
          IF(IPRINT.LE.0)GO TO 80
          WRITE(LP,940)INDEX,CBUFA(I)
940  FORMAT(1X,I5,2X,2E13.4)
80  CONTINUE
C
C BELOW, INVERT ISGN AND IDIR FOR INVERSE TRANSFORM (COMPLEX-TO-REAL)
75  ISGN=-ISGN
      IDIR=-IDIR
      SCAL=1./SIZE
      IF(IRMF.NE.0)
          X CALL RMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX, IPAK)
C EITHER ROUTINE RMFFT INVERSE TRANSFORMS ARRAY TO PACKED REAL
C
      IF(IRMF.EQ.0)
          X CALL CMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX)
C OR ROUTINE CMFFT INVERSE TRANSFORMS ARRAY
C
      IF(IPRINT.LE.0)GO TO 85
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
      WRITE(LP,910)IPRINT,IRMF
      WRITE(LP,920)NDIM,ISGN,IDIR,IBEX,ICEX,IPAK,
      X (MEXA(J),J=1,NDIM)
      WRITE(LP,960)
960  FORMAT(8H INDEX,4X,8HFFT/2**M,6X,5HINPUT,10X,4HDIFF)
C
85  IERR=MFPAR(IRMF,0)
      IF(IERR.NE.0)GO TO 1000
      NCNT=NRBK
      IF(IRMF.EQ.0)NCNT=NCNT/2

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C HELPER ROUTN, NTBK TOTAL NUMB OF NRBK REALS IN FFT BLOCK
C (NCNT IS NUMBER OF ELMTS IN FFT BLOCK, REAL OR CPLX)
C
C BELOW, COMPARES FFT INVERSE RESULT WITH INITIAL RANDOM INPUT
  DIFM=0.
  VALU=0.
  DO 90 JB=1,NTBK
    CALL MFREAD(BUFA,NRBK,JB)
C READ 'RECDS' OF NRBK LENGTH, TOTALLING NTBK (RECOMPUTED BY MFPAR)
C (NOTE, PDP 11 MFREAD/MFWRITE (USING FORTRAN I/O) CANNOT ACCEPT 'RECDS'
C OF DIFFERENT LENGTH; OTHERWISE COULD USE NTD1 'RECDS' OF NR1 HERE)
C
  DO 90 I=1,NCNT
    IF(IRMF.NE.0)RM=BUFA(I)
    IF(IRMF.EQ.0)RM=REAL(CBUFA(I))
    DIF=ABS(RM-VALU)
    IF(DIF.GT.DIFM)DIFM=DIF
    IF(IPRINT.GT.0)WRITE(LP,990)I,RM,VALU,DIF
990  FORMAT(1X,I5,3(2X,E13.4))
    VALU=VALU+VINC
90  CONTINUE
C PRINT INVERSE RESULTS (SHOULD BE SAME AS INITIAL RAMP)
C
  IF(IPRINT.GE.0)WRITE(LP,950)DIFM,NDIM,ISGN,IDIR,IBEX,ICEX,IPAK,
X (MEXA(J),J=1,NDIM)
950  FORMAT(10H MAX DIFF ,E11.4,1X,
X 8H NDIM =,I3,8H ISGN =,I3,8H IDIR =,I3,
X 8H IBEX =,I3,8H ICEX =,I3,8H IPAK =,I3,
X 10H MEXA() =,3I5)
  IF(DIFM.GT.DIFMG)DIFMG=DIFM
C
  STOP
C
1000 WRITE(LP,980)IERR
980  FORMAT(25H IBEX FORCED TOO SMALL OR,
X 30H NOT ALL /MFINT/ CORRECT, IERR,I4)
  STOP
  END

      SUBROUTINE MFREAD(BUFA,NB,JB)
C
C PDP 11 FORTRAN DIRECT ACCESS READ ROUTINE FOR FFT
C
C (LOGICAL UNIT LUN IN COMMON/FFTCOM/ MUST HAVE BEEN OPENED
C PREVIOUSLY; EG. CALL ASSIGN(LUN,'FILENAME',0)
C AND DEFINE FILE LUN(NREC,NWD,U,INDX)
C WHERE NWD=2**(IBEX+1) AND NREC=2**(M-IBEX+1) OR LESS IF IPAK=0 OR -1)
C
C SEE ALSO ALTERNATIVE, FAST MACRO I/O SUBROUTINES
C (THESE ARE TO BE PREFERRED, SINCE SPEED 2 TO 10 TIMES BETTER)
C
C READ BLOCK, INDEX JB, FROM MASS STORE TO BUFA, NB REAL VALUES
C
  COMMON/FFTCOM/LUN
  REAL BUFA(NB)
C
  READ(LUN'JB)BUFA
  RETURN
C
  END

      SUBROUTINE MFWRITE(BUFA,NB,JB)
C
C PDP 11 FORTRAN DIRECT ACCESS WRITE ROUTINE FOR FFT
C
C (LOGICAL UNIT LUN IN COMMON/FFTCOM/ MUST HAVE BEEN OPENED
C PREVIOUSLY; EG. CALL ASSIGN(LUN,'FILENAME',0)
C AND DEFINE FILE LUN(NREC,NWD,U,INDX)
C WHERE NWD=2**(IBEX+1) AND NREC=2**(M-IBEX+1) OR LESS IF IPAK=0 OR -1)
C
C SEE ALSO ALTERNATIVE, FAST MACRO I/O SUBROUTINES
C (THESE ARE TO BE PREFERRED, SINCE SPEED 2 TO 10 TIMES BETTER)
C

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C WRITE BLOCK, INDEX JB, FROM BUFA TO MASS STORE, NB REAL VALUES
C
COMMON/FFTCOM/LUN
REAL BUFA(NB)
C
WRITE(LUN'JB)BUFA
RETURN
C
END

C
C PDP 11 FAST MACRO I/O FFT SAMPLE TEST PROGRAM
C (REQUIRES CALL TO SYSTEM MACRO TO OPEN FILE FOR I/O - SEE LINE 70)
C
C NOTE THAT FAST MACRO I/O IS MORE EFFICIENT THAN STANDARD FORTRAN I/O
C
C NOTE WELL THAT TYPE COMPLEX DATA MUST EXIST AS ALTERNATING
C REAL/IMAG/REAL/IMAG... ELEMENTS, BOTH IN MASS STORE AND IN FORTRAN
C WORKING ARRAY BUFA; IN THIS FFT, DIFFERENT SUBROUTINES WILL SET
C DIFFERENT TYPE (REAL OR COMPLEX) FOR ARRAY BUFA.
C
COMMON/FFTCOM/IFDB,IOERR
C
COMMON /FFTCOM/ HOLDS FDB ADDRESS FOR MACRO I/O (RSX11M), CHANNEL (RT11)
C (SEE COMMENTS IN FAST MACRO ROUTINE MFREAD/MFWRITE)
C
COMMON/MFARG/MEXA(4),NDIM,ISGN,IDIR,SCAL,IBEX,ICEX,IPAK
COMMON/MFVAL/DIMA(4),TDM1,RDM1,FBLK,TBLK,RBLK,RCOR,SIZE
COMMON/MFINT/NDMA(4),NTD1,NRD1,NFBK,NTBK,NRBK,NRCR,NSZE
C
C COMMON AREAS /MFARG/,/MFVAL/,/MFINT/ USEFUL FOR RUNNING MASS STORE FFT
C /MFARG/ HOLDS ARGUMENTS USED IN FFT CALLES, MOSTLY EXPONENTS
C HELPER ROUTINE MFPAR COMPUTES VALUES FROM /MFARG/ INTO
C /MFVAL/ (REALS AS SOME LARGE), /MFINT/ (INTEGER EQUIVALENTS IF POSS)
C OR CAN REVERSE-COMPUTE SOME /MFARG/ FROM /MFVAL/
C SEE COMMENTS, ROUTINE MFPAR.
C
COMPLEX CBUFA(1024)
REAL BUFA(2048)
EQUIVALENCE (CBUFA(1),BUFA(1))
C
C BUFA(1)=CBUFA(1) IS WORKING AREA IN CORE STORE FOR FFT AND PROGRAM
C
LUN=1
C LUN IS LOGICAL UNIT FOR PDP 11 MASS STORE I/O
C
LP=5
IPRINT=0
C LP IS PRINTER LOGICAL UNIT, IPRINT=+1 COPIOUS, 0 MAX DIFFERENCES ONLY
C
IRMF=-1
C IRMF=-1 REAL ROUTINE RMFFT TEST, 0 COMPLEX ROUTINE CMFFT TEST
C
ISGN=-1
IDIR=-1
IPAK=1
C FFT ARGS, ISGN=+/- 1, IDIR=-1 (RMFFT), IDIR=1 OR 0 (CMFFT)
C IPAK=1, 0 OR -1 (RMFFT), NOT USED (CMFFT)
C
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
IF(IPRINT.LE.0)WRITE(LP,910)IPRINT,IRMF
910 FORMAT(31H1MASS STORE FFT TEST - IPRINT =,I3,
X 25H (1=COPIOUS,0=MAX DIFFS),,
X 9H IRMF =,I3,28H (0=CMFFT TEST,1=RMFFT TEST) /)
DIFMG=0.
C
MEXA(1)=8
MEXA(2)=6
NDIM=2
IBEX=7
ICEX=11
C MORE FFT ARGS, 2**8 ROWS OF 2**6 ELMTS, 2 DIMEN,
C FFT MASS STORE BLOCKS 2**IBEX REALS, BUFA 2**ICEX REALS
C

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      IERR=MFPAR(IRMF,0)
      IF(IERR.NE.0)GO TO 1000
C   (NOTE HELPER ROUTN MFPAR RETURNS NFBK AS MAXIMUM FILE SIZE)
C
      STOP 'ASSIGN AND OPEN FILE HERE'
C   DELETE ABOVE LINE AND INVOKE SYSTEM MACRO 'OPEN$' (RSX) OR '.OPEN (RT11)
C
C   FOR EXAMPLE, A FORTRAN CALLABLE SUBROUTINE 'MFOPEN' COULD OPEN FILE
C   'FFTTEST.DAT', RETURNING IFDB=FDB ADDRESS (RSX11M) IN COMMON/FFTCOM/
C   OF CHANNEL NUMBER (RT11), FOR USE BY MACRO MFREAD/MFWRITE, THUS:
C
      CALL MFOPEN(LUN,'FFTTEST.DAT',IFDB)
C
      IOERR=0
C   PRESET IOERR IN /FFTCOM/; ZERO IF NO ERRORS IN I/O
C
      NCNT=NRD1
C   HELPER ROUTINE, NCNT NUM OF ELMTS IN 'FIRST' DIMEN, NTD1 TOTAL BLOCKS
C
      M=MFSUM(MEXA,NDIM,99)
      SCAL=1.
      VALU=0.
      VINC=1./SIZE
      IF(IRMF.EQ.0)GO TO 50
C   SWITCH FOR COMPLEX ROUTINE CMFFT AT 50, REAL RMFFT BELOW
C
      DO 40 JB=1,NTD1
      DO 35 I=1,NCNT
      BUFA(I)=VALU
35     VALU=VALU+VINC
      CALL MFWRITE(BUFA,NRD1,JB)
40     CONTINUE
C   LOAD RAMP FUNCTN IN REAL ARRAY IN 'RECS' OF NRD1 LENGTH
C   (NOTE, 'MACRO' MFREAD/MFWRITE CAN ACCEPT 'RECORDS' OF DIFFERENT LENGTH;
C   SO CAN USE NTD1,NRD1 HERE INSTEAD OF NTBK,NRBK AS IN FORTRAN MASS STORE)
C
      CALL RMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX, IPAK)
C   REAL MASS STORE ROUTINE RMFFT TO TRANSFORM ARRAY TO COMPLEX RESULT
      GO TO 70
C
C   BELOW, USE COMPLEX ROUTINE CMFFT, NCNT NUM OF CPLX ELMTS IN 'FIRST' DIMEN
50     NCNT=NCNT/2
      DO 60 JB=1,NTD1
      DO 55 I=1,NCNT
      CBUFA(I)=CMPLX(VALU,0.)
55     VALU=VALU+VINC
      CALL MFWRITE(BUFA,NRD1,JB)
60     CONTINUE
C   LOAD RAMP FUNCTN IN COMPLEX ARRAY IN 'RECS' OF NRD1 LENGTH
C   (NOTE, 'MACRO' MFREAD/MFWRITE CAN ACCEPT 'RECORDS' OF DIFFERENT LENGTH;
C   SO CAN USE NTD1,NRD1 HERE INSTEAD OF NTBK,NRBK AS IN FORTRAN MASS STORE)
C
      CALL CMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX)
C   COMPLEX MASS STORE ROUTINE CMFFT TO TRANSFORM ARRAY TO COMPLEX RESULT
C
70     IF(IPRINT.LE.0)GO TO 75
C   BELOW, PRINT HEADINGS FOR TEST OUTPUT
      WRITE(LP,910)IPRINT,IRMF
      WRITE(LP,920)NDIM,ISGN,IDIR,IBEX,ICEX,IPAK,
      X (MEXA(J),J=1,NDIM)
920    FORMAT(8H NDIM =,I3,8H ISGN =,I3,8H IDIR =,I3,
      X 8H IBEX =,I3,8H ICEX =,I3,8H IPAK =,I3,
      X 10H MEXA() =,3I5)
      WRITE(LP,930)
930    FORMAT(8H INDEX,12X,3HFFT)
C
      IERR=MFPAR(-IRMF,0)
      IF(IERR.NE.0)GO TO 1000
      NCNT=NRD1/2
C   HELPER ROUTN, NTD1 TOTAL NUMB OF NRD1 REALS IN 'FIRST' DIMEN
C   (NOTE IRMF=+1 FOR DATA IN COMPLEX STATE, NCNT ELMTS)
C
C   BELOW, PROGRAM CAN ACCESS FFT COMPLEX RESULT (AND FORM COMPARISON, IF KNOWN)
      DO 80 JB=1,NTD1
      K=(JB-1)*NCNT
      CALL MFREAD(BUFA,NRD1,JB)

```

```

C READ 'RECS' OF NRD1 LENGTH, TOTALLING NTD1 (RECOMPUTED BY MFPAR)
C (NOTE, 'MACRO' MFREAD/MFWRITE CAN ACCEPT 'RECORDS' OF DIFFERENT LENGTH;
C SO CAN USE NTD1,NRD1 HERE INSTEAD OF NTBK,NRBK AS IN FORTRAN MASS STORE)
C
      DO 80 I=1,NCNT
      INDEX=J-1
      IF(IPRINT.LE.0)GO TO 80
      WRITE(LP,940)INDEX,CBUFA(I)
940  FORMAT(1X,I5,2X,2E13.4)
80   CONTINUE
C
C BELOW, INVERT ISGN AND IDIR FOR INVERSE TRANSFORM (COMPLEX-TO-REAL)
75   ISGN=-ISGN
      IDIR=-IDIR
      SCAL=1./SIZE
      IF(IRMF.NE.0)
        X CALL RMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX, IPAK)
C EITHER ROUTINE RMFFT INVERSE TRANSFORMS ARRAY TO PACKED REAL
C
      IF(IRMF.EQ.0)
        X CALL CMFFT (MEXA,NDIM,ISGN,IDIR,SCAL, BUFA,IBEX,ICEX)
C OR ROUTINE CMFFT INVERSE TRANSFORMS ARRAY
C
      IF(IPRINT.LE.0)GO TO 85
C BELOW, PRINT HEADINGS FOR TEST OUTPUT
      WRITE(LP,910) IPRINT,IRMF
      WRITE(LP,920)NDIM,ISGN,IDIR,IBEX,ICEX,IPAK,
      X (MEXA(J),J=1,NDIM)
      WRITE(LP,960)
960  FORMAT(8H INDEX,4X,8HFFT/2**M,6X,5HINPUT,10X,4HDIFF)
C
85   IERR=MFPAR(IRMF,0)
      IF(IERR.NE.0)GO TO 1000
      NCNT=NRD1
      IF(IRMF.EQ.0)NCNT=NCNT/2
C HELPER ROUTN, NTD1 TOTAL NUMB OF NRD1 REALS IN 'FIRST' DIMEN
C (NCNT IS NUMBER OF ELMTS IN 'FIRST' DIMEN, REAL OR CMLPX)
C
C BELOW, COMPARES FFT INVERSE RESULT WITH INITIAL RANDOM INPUT
      DIFM=0.
      VALU=0.
      DO 90 JB=1,NTD1
      CALL MFREAD(BUFA,NRD1,JB)
C READ 'RECS' OF NRD1 LENGTH, TOTALLING NTD1 (RECOMPUTED BY MFPAR)
C (NOTE, 'MACRO' MFREAD/MFWRITE CAN ACCEPT 'RECORDS' OF DIFFERENT LENGTH;
C SO CAN USE NTD1,NRD1 HERE INSTEAD OF NTBK,NRBK AS IN FORTRAN MASS STORE)
C
      DO 90 I=1,NCNT
      IF(IRMF.NE.0)RM=BUFA(I)
      IF(IRMF.EQ.0)RM=REAL(CBUFA(I))
      DIF=ABS(RM-VALU)
      IF(DIF.GT.DIFM)DIFM=DIF
      IF(IPRINT.GT.0)WRITE(LP,990)I,RM,VALU,DIF
990  FORMAT(1X,I5,3(2X,E13.4))
      VALU=VALU+VINC
90   CONTINUE
C PRINT INVERSE RESULTS (SHOULD BE SAME AS INITIAL RAMP)
C
      IF(IPRINT.GE.0)WRITE(LP,950)DIFM,NDIM,ISGN,IDIR,IBEX,ICEX,IPAK,
      X (MEXA(J),J=1,NDIM)
950  FORMAT(10H MAX DIFF ,E11.4,1X,
      X 8H NDIM =,I3,8H ISGN =,I3,8H IDIR =,I3,
      X 8H IBEX =,I3,8H ICEX =,I3,8H IPAK =,I3,
      X 10H MEXA() =,3I5)
      IF(DIFM.GT.DIFMG)DIFMG=DIFM
C
      STOP
C
1000 WRITE(LP,980)IERR
980  FORMAT(25H IBEX FORCED TOO SMALL OR,
      X 30H NOT ALL /MFINT/ CORRECT, IERR,I4)
      STOP
      END

```

```

;
; PDP11 RSX11M OR RT11 FAST MACRO I/O ROUTINES FOR FFT
;
; SUBROUTINE MFREAD(BUFA,NB,JB)
; SUBROUTINE MFWRIT(BUFA,NB,JB)
;
;C READ BLOCK, INDEX JB, FROM MASS STORE TO BUFA, NB REAL VALUES
;C WRITE BLOCK, INDEX JB, FROM BUFA TO MASS STORE, NB REAL VALUES
;
; COMMON/FFTCOM/IFDB,IOERR (RSX11 FDB ADDRESS, ERROR)
;OR COMMON/FFTCOM/ICHAN,IOERR (RT11 CHANNEL NUM, ERROR)
;
;NOTE1: FILE MUST BE PROPERLY OPEN FOR READ/WRITE ACCESS AND WITH
; RSX11M FDB ADDRESS OR RT11 CHANNEL NUMBER IN FIRST WORD
; OF COMMON/FFTCOM/ (IE. INVOKE RSX11M OR RT11 OPEN MACROS).
;
;NOTE2: I/O ERRORS, IF THEY OCCUR, ARE RETURNED IN SECOND WORD
; OF COMMON/FFTCOM/. THIS WORD (IOERR) REMAINS UNCHANGED
; IF NO ERROR OCCURS - THUS, IF PRESET TO 0 BEFORE CALLING
; MFREAD/MFWRIT (OR MASS STORE FFT), WILL FINALLY BE
; 0 IF NO ERRORS, OR NON ZERO=LAST ERROR (SEE FOLLOWING LABEL L4:).
;
;NOTE3: NB MUST BE AN INTEGRAL POWER OF 2; 128 OR MORE MOST EFFICIENT
; (NB.LE.64 USES LOCAL 256 WORD JBUF TO HOLD SECTOR;
; WRITE ALWAYS WRITES THIS SECTOR TO FILE; SLOW BUT SAFE).
;
;NOTE4: READY TO ASSEMBLE FOR RSX11M. FOR RT11, ERASE FOLLOWING LINE:
RSX=0
;
; .IF DF,RSX
; CONDITIONAL SECTION FOR RSX11M (V3)
; .TITLE MFREAD(RSX11M)
; .MCALL READ$,WRITE$,WAIT$
; .PSECT FFTCOM,RW,D,GBL,REL,OVR
IFDB: .WORD 0 ;FDB ADDRESS
IOERR: .WORD 0 ;ERROR FLAG (UNCHANGED IF NO ERROR)
; .PSECT
; .ENDC
;
; .IF NDF,RSX
; CONDITIONAL SECTION FOR RT11 (V3)
; .TITLE MFREAD(RT11)
; .MCALL .READW,.WRITW
ERRBYT=52
; .PSECT FFTCOM,RW,D,GBL,REL,OVR
ICHAN: .WORD 0 ;CHANNEL NUMBER
IOERR: .WORD 0 ;ERROR FLAG (UNCHANGED IF NO ERROR)
; .PSECT MFREAD,RW,I,LCL,REL,OVR
; .ENDC
;
; .GLOBL MFREAD,MFWRIT
;
MFREAD: MOV #1,R0 ;MFREAD SETS JSW=1
BR L1
;
MFWRIT: MOV #-1,R0 ;MFWRIT SETS JSW=-1
;
L1: MOV R0,JSW ;HOLD JSW
TST (R5)+ ;IGNORE ARGUMENT COUNT
MOV (R5)+,R3 ;R3=ADDRESS OF 'BUFA'
MOV @(R5)+,R2 ;R2='NB'
MOV @(R5)+,R1 ;R1='JB'
; FINISH ACCESSING SUBROUTINE ARGUMENTS
;
; DEC R1 ;R1=JB-1
; ASL R2 ;R2=WORD COUNT, NB*2
; MOV R2,R5 ;COPY TO R5
; BEQ L6 ;FINISH IF COUNT ZERO
; CLR R0 ;CLEAR 0 FOR WORD OFFSET
; CMP R5,#256. ;CHECK WHETHER WHOLE SECTORS
; BLT L12 ;GO TO L12 IF NOT (MORE COMPLEX)
;
; BELOW, SCALE R1 TO SECTOR OFFSET (ASSUMES NWD POWER OF 2)
L10: BIT #256.,R5 ;LOOK FOR BIT AT 256.
BNE L11 ;R1 COMPLETE WHEN FOUND
ASL R1 ;SCALE R1 ACCORDING TO R5

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        ASR      R5
        BR      L1Ø
;
; NOW HAVE    R1=SECTOR START (FIRST=Ø)
;            R2=WORD COUNT
;            R3=BUFA ADDRESS
;
L11:   MOV      JSW,JSWIO      ;SET UP FOR READ/WRITE
        BGT      L13
        MOV      #-1,SECHLD    ;DIRECT WRITE, ENSURE JBUF LOOKS EMPTY
L13:   JSR      PC,INOUT      ;CALL DIRECT INPUT/CUTPUT
L6:    RTS      PC            ;RETURN FROM SUBROUTINE
; END OF SIMPLE DIRECT INPUT/OUTPUT
;
        .IF     DF,RSX
; CONDITIONAL SECTION FOR RSX11M
INOUT: ASL      R2            ;RSX11M NEEDS BYTE COUNT=NB*4
        INC      R1            ;RSX11M SECTOR STARTS WITH 1
        MOV      R1,SECTL     ;HOLD AS LOWER SECTOR VALUE
        MOV      IFDB,R4      ;R4=FDB ADDRESS
        MOVB     F.LUN(R4),R1 ;USE LUN FOR EVENT FLAG
;
        TST      JSWIO
        BLT      L5            ;BRANCH IF WRITE
;
        MOV      F.EFBK+2(R4),RØ ;GET LOW SECTOR OF EOF
        DEC      RØ            ;ALLOW FOR HEADER SECTOR
        CMP      RØ,SECTL     ;TEST IF SECTOR EXISTS YET
        BLT      L7            ;NO READ IF DOES NOT EXIST
;
        READ$    R4,R3,R2,#SECT,R1 ;FDB,BUF,BYTCNT,SECT,EVFLG
;
L4:    MOV      IFDB,R4      ;R4=FDB ADDRESS
        MOVB     F.LUN(R4),R1 ;USE LUN FOR EVENT FLAG
        WAIT$    R4,R1      ;WAIT FOR I/O FINISH
        BCC      L7            ;NO ERROR IF CARRY CLEAR
;
        MOV      IFDB,R4      ;R4=FDB ADDRESS
        MOVB     F.ERR(R4),R4 ;HOLD NEGATIVE ERROR CODE
        MOV      R4,IOERR     ;IN IOERR IN COMMON/FFTCOM/
L7:    RTS      PC
;
L5:    WRITE$    R4,R3,R2,#SECT,R1 ;FDB,BUF,BYTCNT,SECT,EVFLG
        BR      L4
;
SECT:  .WORD    Ø            ;HOLDS SECTOR VALUE (HIGHER, ALWAYS Ø)
SECTL: .WORD    Ø            ;(LOWER, COMPUTED)
        .ENDC
;
        .IF     NDF,RSX
;CONDITIONAL SECTION FOR RT11
INOUT: MOV      ICHAN,R4      ;R4=CHANNEL NUMBER
;
        TST      JSWIO
        BLT      L5            ;BRANCH IF WRITE
;
        .READW   #AREA,R4,R3,R2,R1 ;AREA,CHAN,BUF,WDCNT,SECT
;
L4:    BCC      L7            ;NO ERROR IF CARRY CLEAR
;
        MOVB     ERRBYT,R4    ;HOLD ERROR CODE + 1 (TO MAKE NON ZERO)
        INC      R4            ;ERRORS Ø,1,2 GO TO 1,2,3
        MOV      R4,IOERR     ;IN IOERR IN COMMON/FFTCOM/
L7:    RTS      PC
;
L5:    .WRITEW   #AREA,R4,R3,R2,R1 ;AREA,CHAN,BUF,WDCNT,SECT
        BR      L4
AREA:  .BLKW    1Ø          ;AREA FOR RT11 I/O MACROS USE
        .ENDC
;
; BELOW, MORE COMPLEX HANDLING OF NWD.LT.256 (PART SECTORS)
L12:   CLC
        ROR      R1            ;SCALE R1/RØ RIGHT ACCORDING TO R5
        ROR      RØ            ;RØ HOLDS FLOW OUT OF R1
        ASL      R5
        BIT      #256.,R5     ;LOOK FOR BIT 256

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        BNE     L20          ;R1/R0 COMPLETE IF FOUND
        BR      L12
;
L20:    SWAB     R0
        ASL     R0
;
; NOW HAVE     R0=BYTE OFFSET IN LOCAL BUFFER
;             R1=SECTOR REQUIRED
;             R2=WORD COUNT
;             R3=USER BUFA ADDRESS
;
; BELOW, TRANSFER BETWEEN LOCAL AND USER BUFFER
        ADD     #JBUF,R0      ;R0=JBUF ADDRESS + OFFSET
        CMP     SECHLD,R1     ;CHECK CURRENT SECTOR LOADED
        BNE     L24          ;MUST FIRST LOAD SECTO AT L24 IF NEC
L21:    TST     JSW          ;TEST WHETHER READ/WRITE
        BLT     L23
;
; BELOW, 'READ' PART SECTOR
L22:    MOV     (R0)+,(R3)+   ;COPY LOCAL TO USER BUFA
        DEC     R2
        BGT     L22
        JMP     L6           ;FINISH
;
; BELOW, 'WRITE' PART SECTOR
L23:    MOV     (R3)+,(R0)+   ;COPY USER TO LOCAL BUF
        DEC     R2
        BGT     L23
        MOV     #-1,JSWIO     ;SET UP I/O FOR WRITE
        JSR     PC,LINOUT     ;WRITE FROM LOCAL JBUF (SLOW BUT SAFE)
        JMP     L6           ;FINISH
;
; BELOW, LOAD LOCAL JBUF, IF POSSIBLE
L24:    MOV     R1,SECHLD     ;HOLD CURRENT SECTOR NUM
        MOV     #1,JSWIO     ;SET UP FOR READ
        JSR     PC,LINOUT     ;DO LOCAL READ TO JBUF
        BR      L21
;
; BELOW, LOCAL INPUT/OUTPUT ROUTINE
LINOUT: MOV     R0,-(SP)      ;SAVE REGISTERS R0 TO R3
        MOV     R1,-(SP)
        MOV     R2,-(SP)
        MOV     R3,-(SP)
        MOV     #256.,R2     ;R2=256. FOR WHOLE SECTOR
        MOV     #JBUF,R3     ;R3 IS LOCAL JBUF ADDRESS
        JSR     PC,INOUT     ;CALL INOUT ROUTINE
        MOV     (SP)+,R3     ;RESTORE REGISTERS R0 TO R3
        MOV     (SP)+,R2
        MOV     (SP)+,R1
        MOV     (SP)+,R0
        RTS     PC
;
; BELOW, LOCAL VARIABLES
JSW:    .WORD   0             ;JSW=+1 MFREAD, -1 MFWRIT
JSWIO:  .WORD   0             ;ACTUAL JSW USED FOR INOUT ROUTINE
SECHLD: .WORD   -1          ;CURRENT SECTOR LOADED IN JBUF
JBUF:   .BLKW  256.         ;JBUF 256 WORD LOCAL BUFFER FOR NWD.LT.256
;
        .END

```


ALGORITHM 546

SOLVEBLOK [F4]

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Key Words and Phrases: almost block diagonal systems, Gaussian elimination, spline approximation, ordinary differential equations

CR Categories: 5.13, 5.14, 5.17

Language: Fortran

DESCRIPTION

This Fortran package is a complement to [1] where its design is explained and it is compared with related algorithms.

REFERENCES

1. DE BOOR, C., AND WEISS, R. SOLVEBLOK: A package for solving almost block diagonal linear systems. *ACM Trans. Math. Software* 6, 1 (March 1980), 80-87.

ALGORITHM

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      SUBROUTINE SLVBLK ( BLOKS, INTEGS, NBLOKS, B, IPIVOT, X, IFLAG ) BLK00100
C   THIS PROGRAM SOLVES THE LINEAR SYSTEM A*X = B WHERE A IS AN BLK00200
C   ALMOST BLOCK DIAGONAL MATRIX. SUCH ALMOST BLOCK DIAGONAL MATRICES BLK00300
C   ARISE NATURALLY IN PIECEWISE POLYNOMIAL INTERPOLATION OR APPROX- BLK00400
C   IMATION AND IN FINITE ELEMENT METHODS FOR TWO-POINT BOUNDARY VALUE BLK00500
C   PROBLEMS. THE PLU FACTORIZATION METHOD IS IMPLEMENTED HERE TO TAKE BLK00600
C   ADVANTAGE OF THE SPECIAL STRUCTURE OF SUCH SYSTEMS FOR SAVINGS IN BLK00700
C   COMPUTING TIME AND STORAGE REQUIREMENTS. BLK00800
C   BLK00900
C   PARAMETERS BLK01000
C   BLOKS  A ONE-DIMENSIONAL ARRAY, OF LENGTH BLK01100
C           SUM( INTEGS(1,I)*INTEGS(2,I) , I = 1,NBLOKS ) BLK01200
C   ON INPUT, CONTAINS THE BLOCKS OF THE ALMOST BLOCK DIAGONAL BLK01300
C   MATRIX A . THE ARRAY INTEGS (SEE BELOW AND THE EXAMPLE) BLK01400
C   DESCRIBES THE BLOCK STRUCTURE. BLK01500
C   ON OUTPUT, CONTAINS CORRESPONDINGLY THE PLU FACTORIZATION BLK01600
C   OF A (IF IFLAG .NE. 0). CERTAIN OF THE ENTRIES INTO BLOKS BLK01700
C   ARE ARBITRARY (WHERE THE BLOCKS OVERLAP). BLK01800
C   INTEGS INTEGER ARRAY DESCRIPTION OF THE BLOCK STRUCTURE OF A . BLK01900
C           INTEGS(1,I) = NO. OF ROWS OF BLOCK I           = NROW BLK02000
C           INTEGS(2,I) = NO. OF COLUMNS OF BLOCK I       = NCOL BLK02100
C           INTEGS(3,I) = NO. OF ELIM. STEPS IN BLOCK I    = LAST BLK02200

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C             I = 1,2,...,NBLOKS                BLK02300
C     THE LINEAR SYSTEM IS OF ORDER            BLK02400
C     N = SUM ( INTEGS(3,I) , I=1,...,NBLOKS ), BLK02500
C     BUT THE TOTAL NUMBER OF ROWS IN THE BLOCKS IS BLK02600
C     NBROWS = SUM( INTEGS(1,I) , I = 1,...,NBLOKS) BLK02700
C     NBLOKS NUMBER OF BLOCKS                BLK02800
C     B RIGHT SIDE OF THE LINEAR SYSTEM, ARRAY OF LENGTH NBROWS. BLK02900
C     CERTAIN OF THE ENTRIES ARE ARBITRARY, CORRESPONDING TO BLK03000
C     ROWS OF THE BLOCKS WHICH OVERLAP (SEE BLOCK STRUCTURE AND BLK03100
C     THE EXAMPLE BELOW).                  BLK03200
C     IPIVOT ON OUTPUT, INTEGER ARRAY CONTAINING THE PIVOTING SEQUENCE BLK03300
C     USED. LENGTH IS NBROWS                BLK03400
C     X ON OUTPUT, CONTAINS THE COMPUTED SOLUTION (IF IFLAG .NE. 0) BLK03500
C     LENGTH IS N.                          BLK03600
C     IFLAG ON OUTPUT, INTEGER              BLK03700
C     = (-1)**(NO. OF INTERCHANGES DURING FACTORIZATION) BLK03800
C     IF A IS INVERTIBLE                    BLK03900
C     = 0 IF A IS SINGULAR                  BLK04000
C
C     AUXILIARY PROGRAMS                    BLK04100
C     FCBLK (BLOKS,INTEGS,NBLOKS,IPIVOT,SCRATCH,IFLAG) FACTORS THE MATRIX BLK04200
C     A , AND IS USED FOR THIS PURPOSE IN SLVBLK. ITS ARGUMENTS BLK04300
C     ARE AS IN SLVBLK, EXCEPT FOR       BLK04400
C     SCRATCH = A WORK ARRAY OF LENGTH MAX(INTEGS(1,I)). BLK04500
C
C     SBBLOK (BLOKS,INTEGS,NBLOKS,IPIVOT,B,X) SOLVES THE SYSTEM A*X = B BLK04600
C     ONCE A IS FACTORED. THIS IS DONE AUTOMATICALLY BY SLVBLK BLK04700
C     FOR ONE RIGHT SIDE B, BUT SUBSEQUENT SOLUTIONS MAY BE BLK04800
C     OBTAINED FOR ADDITIONAL B-VECTORS. THE ARGUMENTS ARE ALL BLK04900
C     AS IN SLVBLK.                        BLK05000
C
C     DTBLOK (BLOKS,INTEGS,NBLOKS,IPIVOT,IFLAG,DETSGN,DETLOG) COMPUTES THE BLK05100
C     DETERMINANT OF A ONCE SLVBLK OR FCBLK HAS DONE THE FACT- BLK05200
C     ORIZATION.THE FIRST FIVE ARGUMENTS ARE AS IN SLVBLK. BLK05300
C     DETSGN = SIGN OF THE DETERMINANT     BLK05400
C     DETLOG = NATURAL LOG OF THE DETERMINANT BLK05500
C
C     ----- BLOCK STRUCTURE OF A ----- BLK05600
C     THE NBLOKS BLOCKS ARE STORED CONSECUTIVELY IN THE ARRAY BLOKS . BLK05700
C     THE FIRST BLOCK HAS ITS (1,1)-ENTRY AT BLOKS(1), AND, IF THE I-TH BLK05800
C     BLOCK HAS ITS (1,1)-ENTRY AT BLOKS(INDEX(I)), THEN BLK05900
C     INDEX(I+1) = INDEX(I) + NROW(I)*NCOL(I) . BLK06000
C     THE BLOCKS ARE PIECED TOGETHER TO GIVE THE INTERESTING PART OF A BLK06100
C     AS FOLLOWS. FOR I = 1,2,...,NBLOKS-1, THE (1,1)-ENTRY OF THE NEXT BLK06200
C     BLOCK (THE (I+1)ST BLOCK ) CORRESPONDS TO THE (LAST+1, LAST+1)-ENTRY BLK06300
C     OF THE CURRENT I-TH BLOCK. RECALL LAST = INTEGS(3,I) AND NOTE THAT BLK06400
C     THIS MEANS THAT                      BLK06500
C     A. EVERY BLOCK STARTS ON THE DIAGONAL OF A . BLK06600
C     B. THE BLOCKS OVERLAP (USUALLY). THE ROWS OF THE (I+1)ST BLOCK BLK06700
C     WHICH ARE OVERLAPPED BY THE I-TH BLOCK MAY BE ARBITRARILY DE- BLK06800
C     FINED INITIALLY. THEY ARE OVERWRITTEN DURING ELIMINATION. BLK06900
C     THE RIGHT SIDE FOR THE EQUATIONS IN THE I-TH BLOCK ARE STORED COR- BLK07000
C     RESPONDINGLY AS THE LAST ENTRIES OF A PIECE OF B OF LENGTH NROW BLK07100
C     (= INTEGS(1,I)) AND FOLLOWING IMMEDIATELY IN B THE CORRESPONDING BLK07200
C     PIECE FOR THE RIGHT SIDE OF THE PRECEDING BLOCK, WITH THE RIGHT SIDE BLK07300
C     FOR THE FIRST BLOCK STARTING AT B(1) . IN THIS, THE RIGHT SIDE FOR BLK07400
C     AN EQUATION NEED ONLY BE SPECIFIED ONCE ON INPUT, IN THE FIRST BLOCK BLK07500
C     IN WHICH THE EQUATION APPEARS.      BLK07600
C
C     ----- EXAMPLE AND TEST DRIVER ----- BLK07700
C     THE TEST DRIVER FOR THIS PACKAGE CONTAINS AN EXAMPLE, A LINEAR BLK07800
C     SYSTEM OF ORDER 11, WHOSE NONZERO ENTRIES ARE INDICATED IN THE FOL- BLK07900
C     LOWING SCHEMA BY THEIR ROW AND COLUMN INDEX MODULO 10. NEXT TO IT BLK08000
C     ARE THE CONTENTS OF THE INTEGS ARRAY WHEN THE MATRIX IS TAKEN TO BLK08100
C     BE ALMOST BLOCK DIAGONAL WITH NBLOKS = 5, AND BELOW IT ARE THE FIVE BLK08200
C     BLOCKS.                               BLK08300
C
C     NROW1 = 3, NCOL1 = 4                  BLK08400
C     11 12 13 14                          BLK08500
C     21 22 23 24  NROW2 = 3, NCOL2 = 3    BLK08600
C     31 32 33 34                          BLK08700
C     LAST1 = 2 43 44 45                   BLK08800
C     53 54 55                             BLK08900
C     NROW3 = 3, NCOL3 = 4                 BLK09000
C     LAST2 = 3 66 67 68 69  NROW4 = 3, NCOL4 = 4 BLK09100
C     76 77 78 79  NROW5 = 4, NCOL5 = 4    BLK09200
C     86 87 88 89                          BLK09300
C     LAST3 = 1 97 98 99 90                BLK09400
C

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```

C          LAST4 = 1   08 09 00 01          BLK10000
C          18 19 10 11          BLK10100
C          LAST5 = 4          BLK10200
C          BLK10300
C          ACTUAL INPUT TO BLOKS SHOWN BY ROWS OF BLOCKS OF A .          BLK10400
C          (THE ** ITEMS ARE ARBITRARY, THIS STORAGE IS USED BY SLVBK)          BLK10500
C          BLK10600
C 11 12 13 14 / ** ** ** / 66 67 68 69 / ** ** ** ** / ** ** ** **          BLK10700
C 21 22 23 24 / 43 44 45 / 76 77 78 79 / ** ** ** ** / ** ** ** *          BLK10800
C 31 32 33 34/ 53 54 55/ 86 87 88 89/ 97 98 99 90/ 08 09 00 01          BLK10900
C          18 19 10 11          BLK11000
C          BLK11100
C INDEX = 1      INDEX = 13  INDEX = 22      INDEX = 34      INDEX = 46          BLK11200
C          BLK11300
C          ACTUAL RIGHT SIDE VALUES WITH ** FOR ARBITRARY VALUES          BLK11400
C B1 B2 B3 ** B4 B5 B6 B7 B8 ** ** B9 ** ** B10 B11          BLK11500
C          BLK11600
C (IT WOULD HAVE BEEN MORE EFFICIENT TO COMBINE BLOCK 3 WITH BLOCK 4)          BLK11700
C          BLK11800
C          INTEGER INTEG(3,NBLOKS),IPIVOT(1),IFLAG          BLK11900
C          REAL BLOKS(1),B(1),X(1)          BLK12000
C          IN THE CALL TO FCBLOK, X IS USED FOR TEMPORARY STORAGE.          BLK12100
C          CALL FCBLOK(BLOKS,INTEGS,NBLOKS,IPIVOT,X,IFLAG)          BLK12200
C          IF (IFLAG .EQ. 0)          RETURN          BLK12300
C          CALL SBBLOK(BLOKS,INTEGS,NBLOKS,IPIVOT,B,X)          BLK12400
C          RETURN          BLK12500
C          END          BLK12600

SUBROUTINE FCBLOK ( BLOKS, INTEG, NBLOKS, IPIVOT, SCRTCH, IFLAG )          BLK12700
CALLS SUBROUTINES F A C T R B AND S H I F T B .          BLK12800
C          BLK12900
C F C B L O K SUPERVISES THE PLU FACTORIZATION WITH PIVOTING OF          BLK13000
C SCALED ROWS OF THE ALMOST BLOCK DIAGONAL MATRIX STORED IN THE ARRAYS          BLK13100
C B L O K S AND I N T E G S .          BLK13200
C          BLK13300
C FACTRB = SUBPROGRAM WHICH CARRIES OUT STEPS 1,...,LAST OF GAUSS          BLK13400
C ELIMINATION (WITH PIVOTING) FOR AN INDIVIDUAL BLOCK.          BLK13500
C SHIFTB = SUBPROGRAM WHICH SHIFTS THE REMAINING ROWS TO THE TOP OF          BLK13600
C THE NEXT BLOCK          BLK13700
C          BLK13800
C          BLK13900
C PARAMETERS          BLK14000
C BLOKS AN ARRAY THAT INITIALLY CONTAINS THE ALMOST BLOCK DIAGONAL          BLK14100
C MATRIX A TO BE FACTORED, AND ON RETURN CONTAINS THE COM-          BLK14200
C PUTED FACTORIZATION OF A .          BLK14300
C INTEG AN INTEGER ARRAY DESCRIBING THE BLOCK STRUCTURE OF A .          BLK14400
C NBLOKS THE NUMBER OF BLOCKS IN A .          BLK14500
C IPIVOT AN INTEGER ARRAY OF DIMENSION SUM (INTEGS(1,I) , I=1,          BLK14600
C ...,NBLOKS) WHICH, ON RETURN, CONTAINS THE PIVOTING STRA-          BLK14700
C TEGY USED.          BLK14800
C SCRTCH WORK AREA REQUIRED, OF LENGTH MAX (INTEGS(1,I) , I=1,          BLK14900
C ...,NBLOKS).          BLK15000
C IFLAG OUTPUT PARAMETER,          BLK15100
C = 0 IN CASE MATRIX WAS FOUND TO BE SINGULAR.          BLK15200
C OTHERWISE,          BLK15300
C = (-1)**(NUMBER OF ROW INTERCHANGES DURING FACTORIZATION)          BLK15400
C          BLK15500
C          INTEGER INTEG(3,NBLOKS),IPIVOT(1),IFLAG, I,INDEX,INDEXB,INDEXN,          BLK15600
C          * LAST,NCOL,NROW          BLK15700
C          REAL BLOKS(1),SCRTCH(1)          BLK15800
C          IFLAG = 1          BLK15900
C          INDEXB = 1          BLK16000
C          INDEXN = 1          BLK16100
C          I = 1          BLK16200
C          LOOP OVER THE BLOCKS. I IS LOOP INDEX          BLK16300
10 INDEX = INDEXN          BLK16400
C          NROW = INTEG(1,I)          BLK16500
C          NCOL = INTEG(2,I)          BLK16600
C          LAST = INTEG(3,I)          BLK16700
C          CARRY OUT ELIMINATION ON THE I-TH BLOCK UNTIL NEXT BLOCK          BLK16800
C          ENTERS, I.E., FOR COLUMNS 1,...,LAST OF I-TH BLOCK.          BLK16900
C          CALL FACTRB(BLOKS(INDEX),IPIVOT(INDEXB),SCRTCH,NROW,NCOL,LAST,          BLK17000
C          * IFLAG)          BLK17100
C          CHECK FOR HAVING REACHED A SINGULAR BLOCK OR THE LAST BLOCK          BLK17200
C          IF (IFLAG .EQ. 0 .OR. I .EQ. NBLOKS)          BLK17300
C          *          RETURN

```

```

      I = I+1                                BLK17400
      INDEXN = NROW*NCOL + INDEX             BLK17500
C      PUT THE REST OF THE I-TH BLOCK ONTO THE NEXT BLOCK BLK17600
      CALL SHIFTB(BLOKS(INDEX),IPIVOT(INDEXB),NROW,NCOL,LAST,
*      BLOKS(INDEXN),INTEGS(1,I),INTEGS(2,I)) BLK17700
      INDEXB = INDEXB + NROW                 BLK17800
      GO TO 10                               BLK17900
      END                                    BLK18000
      BLK18100

      SUBROUTINE FACTRB ( W, IPIVOT, D, NROW, NCOL, LAST, IFLAG ) BLK18200
C ADAPTED FROM P.132 OF *ELEMENTARY NUMER.ANALYSIS* BY CONTE-DE BOOR BLK18300
C                                                                    BLK18400
C CONSTRUCTS A PARTIAL PLU FACTORIZATION, CORRESPONDING TO STEPS 1,...,BLK18500
C L A S T IN GAUSS ELIMINATION, FOR THE MATRIX W OF ORDER BLK18600
C ( N R O W , N C O L ), USING PIVOTING OF SCALED ROWS. BLK18700
C                                                                    BLK18800
C PARAMETERS BLK18900
C W CONTAINS THE (NROW,NCOL) MATRIX TO BE PARTIALLY FACTORED BLK19000
C ON INPUT, AND THE PARTIAL FACTORIZATION ON OUTPUT. BLK19100
C IPIVOT AN INTEGER ARRAY OF LENGTH NROW CONTAINING A RECORD OF THE BLK19200
C PIVOTING STRATEGY USED. ROW IPIVOT(I) IS USED DURING THE BLK19300
C I-TH ELIMINATION STEP, I=1,...,LAST. BLK19400
C D A WORK ARRAY OF LENGTH NROW USED TO STORE ROW SIZES BLK19500
C TEMPORARILY. BLK19600
C NROW NUMBER OF ROWS OF W. BLK19700
C NCOL NUMBER OF COLUMNS OF W. BLK19800
C LAST NUMBER OF ELIMINATION STEPS TO BE CARRIED OUT. BLK19900

C IFLAG ON OUTPUT, EQUALS IFLAG ON INPUT TIMES (-1)**(NUMBER OF BLK20000
C ROW INTERCHANGES DURING THE FACTORIZATION PROCESS), IN BLK20100
C CASE NO ZERO PIVOT WAS ENCOUNTERED. BLK20200
C OTHERWISE, IFLAG = 0 ON OUTPUT. BLK20300
C BLK20400
C INTEGER IPIVOT(NROW),NCOL,LAST,IFLAG, I,IPIVI,IPIVK,J,K,KP1 BLK20500
C REAL W(NROW,NCOL),D(NROW), AWIKDI,COLMAX,RATIO,ROWMAX BLK20600
C INITIALIZE IPIVOT, D BLK20700
  DO 10 I=1,NROW BLK20800
    IPIVOT(I) = I BLK20900
    ROWMAX = 0. BLK21000
    DO 9 J=1,NCOL BLK21100
      ROWMAX = AMAX1(ROWMAX, ABS(W(I,J))) BLK21200
    IF (ROWMAX .EQ. 0.) GO TO 999 BLK21300
  10 D(I) = ROWMAX BLK21400
C GAUSS ELIMINATION WITH PIVOTING OF SCALED ROWS, LOOP OVER K=1,..,LAST BLK21500
  K = 1 BLK21600
C AS PIVOT ROW FOR K-TH STEP, PICK AMONG THE ROWS NOT YET USED, BLK21700
C I.E., FROM ROWS IPIVOT(K),...,IPIVOT(NROW), THE ONE WHOSE K-TH BLK21800
C ENTRY (COMPARED TO THE ROW SIZE) IS LARGEST. THEN, IF THIS ROW BLK21900
C DOES NOT TURN OUT TO BE ROW IPIVOT(K), REDEFINE IPIVOT(K) AP- BLK22000
C PROPRIATELY AND RECORD THIS INTERCHANGE BY CHANGING THE SIGN BLK22100
C OF I F L A G . BLK22200
  11 IPIVK = IPIVOT(K) BLK22300
  IF (K .EQ. NROW) GO TO 21 BLK22400
  J = K BLK22500
  KP1 = K+1 BLK22600
  COLMAX = ABS(W(IPIVK,K))/D(IPIVK) BLK22700
C FIND THE (RELATIVELY) LARGEST PIVOT BLK22800
  DO 15 I=KP1,NROW BLK22900
    IPIVI = IPIVOT(I) BLK23000
    AWIKDI = ABS(W(IPIVI,K))/D(IPIVI) BLK23100
    IF (AWIKDI .LE. COLMAX) GO TO 15 BLK23200
    COLMAX = AWIKDI BLK23300
    J = I BLK23400
  15 CONTINUE BLK23500
  IF (J .EQ. K) GO TO 16 BLK23600
  IPIVK = IPIVOT(J) BLK23700
  IPIVOT(J) = IPIVOT(K) BLK23800
  IPIVOT(K) = IPIVK BLK23900
  IFLAG = -IFLAG BLK24000
  16 CONTINUE BLK24100
C IF PIVOT ELEMENT IS TOO SMALL IN ABSOLUTE VALUE, DECLARE BLK24200
C MATRIX TO BE NONINVERTIBLE AND QUIT. BLK24300
  IF (ABS(W(IPIVK,K))+D(IPIVK) .LE. D(IPIVK)) BLK24400
* GO TO 999 BLK24500
C OTHERWISE, SUBTRACT THE APPROPRIATE MULTIPLE OF THE PIVOT BLK24600

```



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SUBROUTINE SBBLOK ( BLOKS, INTEGS, NBLOKS, IPIVOT, B, X )      BLK31700
CALLS SUBROUTINES S U B F O R AND S U B B A K .                BLK31800
C                                                                BLK31900
C SUPERVISES THE SOLUTION (BY FORWARD AND BACKWARD SUBSTITUTION) OF BLK32000
C THE LINEAR SYSTEM A*X = B FOR X, WITH THE PLU FACTORIZATION OF A BLK32100
C ALREADY GENERATED IN F C B L O K . INDIVIDUAL BLOCKS OF EQUATIONS BLK32200
C ARE SOLVED VIA S U B F O R AND S U B B A K .                  BLK32300
C                                                                BLK32400
C PARAMETERS                                                    BLK32500
C   BLOKS, INTEGS, NBLOKS, IPIVOT ARE AS ON RETURN FROM FCBLK. BLK32600
C   B THE RIGHT SIDE, STORED CORRESPONDING TO THE STORAGE OF BLK32700
C   THE EQUATIONS. SEE COMMENTS IN S L V B L K FOR DETAILS. BLK32800
C   X SOLUTION VECTOR                                           BLK32900
C                                                                BLK33000
C   INTEGER INTEGS(3,NBLOKS),IPIVOT(1), I,INDEX,INDEXB,INDEXX,J,LAST, BLK33100
C   * NBP1,NCOL,NROW                                           BLK33200
C   REAL BLOKS(1),B(1),X(1)                                     BLK33300
C                                                                BLK33400
C   FORWARD SUBSTITUTION PASS                                  BLK33500
C                                                                BLK33600
C   INDEX = 1                                                  BLK33700
C   INDEXB = 1                                                 BLK33800
C   INDEXX = 1                                                 BLK33900
C   DO 20 I=1,NBLOKS                                          BLK34000
C     NROW = INTEGS(1,I)                                       BLK34100
C     LAST = INTEGS(3,I)                                       BLK34200
C     CALL SUBFOR(BLOKS(INDEX),IPIVOT(INDEXB),NROW,LAST,B(INDEXB), BLK34300
C     * X(INDEXX))                                              BLK34400
C     INDEX = NROW*INTEGS(2,I) + INDEX                         BLK34500
C     INDEXB = INDEXB + NROW                                   BLK34600
C 20  INDEXX = INDEXX + LAST                                   BLK34700
C                                                                BLK34800
C   BACK SUBSTITUTION PASS                                    BLK34900
C                                                                BLK35000
C   NBP1 = NBLOKS + 1                                         BLK35100
C   DO 30 J=1,NBLOKS                                          BLK35200
C     I = NBP1 - J                                             BLK35300
C     NROW = INTEGS(1,I)                                       BLK35400
C     NCOL = INTEGS(2,I)                                       BLK35500
C     LAST = INTEGS(3,I)                                       BLK35600
C     INDEX = INDEX - NROW*NCOL                                BLK35700
C     INDEXB = INDEXB - NROW                                   BLK35800
C     INDEXX = INDEXX - LAST                                   BLK35900
C 30  CALL SUBBAK(BLOKS(INDEX),IPIVOT(INDEXB),NROW,NCOL,LAST, BLK36000
C   * X(INDEXX))                                              BLK36100
C                                                                BLK36200
C   RETURN                                                    BLK36300
C   END

```

```

SUBROUTINE SUBFOR ( W, IPIVOT, NROW, LAST, B, X )              BLK36400
C CARRIES OUT THE FORWARD PASS OF SUBSTITUTION FOR THE CURRENT BLOCK, BLK36500
C I.E., THE ACTION ON THE RIGHT SIDE CORRESPONDING TO THE ELIMINATION BLK36600
C CARRIED OUT IN F A C T R B FOR THIS BLOCK.                  BLK36700
C AT THE END, X(J) CONTAINS THE RIGHT SIDE OF THE TRANSFORMED BLK36800
C IPIVOT(J)-TH EQUATION IN THIS BLOCK, J=1,...,NROW. THEN, SINCE BLK36900
C FOR I=1,...,NROW-LAST, B(NROW+I) IS GOING TO BE USED AS THE RIGHT BLK37000
C SIDE OF EQUATION I IN THE NEXT BLOCK (SHIFTED OVER THERE FROM BLK37100
C THIS BLOCK DURING FACTORIZATION), IT IS SET EQUAL TO X(LAST+I) HERE. BLK37200
C                                                                BLK37300
C PARAMETERS                                                    BLK37400
C   W, IPIVOT, NROW, LAST ARE AS ON RETURN FROM FACTRB.      BLK37500
C   B(J) IS EXPECTED TO CONTAIN, ON INPUT, THE RIGHT SIDE OF J-TH BLK37600
C EQUATION FOR THIS BLOCK, J=1,...,NROW.                     BLK37700
C   B(NROW+J) CONTAINS, ON OUTPUT, THE APPROPRIATELY MODIFIED RIGHT BLK37800
C SIDE FOR EQUATION J IN NEXT BLOCK, J=1,...,NROW-LAST.     BLK37900
C   X(J) CONTAINS, ON OUTPUT, THE APPROPRIATELY MODIFIED RIGHT BLK38000
C SIDE OF EQUATION IPIVOT(J) IN THIS BLOCK, J=1,...,LAST (AND BLK38100
C EVEN FOR J=LAST+1,...,NROW).                                BLK38200
C                                                                BLK38300
C   INTEGER IPIVOT(NROW), IP,JMAX,K                           BLK38400
C   DIMENSION B(NROW + NROW-LAST)                             BLK38500
C   REAL W(NROW,LAST),B(1),X(NROW)                            BLK38600
C   IP = IPIVOT(1)                                            BLK38700
C   X(1) = B(IP)                                              BLK38800
C   IF (NROW.EQ. 1)                                           BLK38900
C     GO TO 99
C   DO 15 K=2,NROW                                           BLK39000

```

```

        IP = IPIVOT(K)
        JMAX = AMIN0(K-1, LAST)
        SUM = 0.
        DO 14 J=1, JMAX
14         SUM = W(IP, J)*X(J) + SUM
15         X(K) = B(IP) - SUM
C
C     TRANSFER MODIFIED RIGHT SIDES OF EQUATIONS IPIVOT(LAST+1), ...,
C     IPIVOT(NROW) TO NEXT BLOCK.
        NROWML = NROW - LAST
        IF (NROWML .EQ. 0)                GO TO 99
        LASTP1 = LAST+1
        DO 25 K=LASTP1, NROW
25         B(NROWML+K) = X(K)
99         RETURN
        END
        BLK39100
        BLK39200
        BLK39300
        BLK39400
        BLK39500
        BLK39600
        BLK39700
        BLK39800
        BLK39900
        BLK40000
        BLK40100
        BLK40200
        BLK40300
        BLK40400
        BLK40500
        BLK40600

        SUBROUTINE SUBBAK ( W, IPIVOT, NROW, NCOL, LAST, X )
C     CARRIES OUT BACKSUBSTITUTION FOR CURRENT BLOCK.
C
C     PARAMETERS
C     W, IPIVOT, NROW, NCOL, LAST ARE AS ON RETURN FROM FACTRB.
C     X(1), ..., X(NCOL) CONTAINS, ON INPUT, THE RIGHT SIDE FOR THE
C     EQUATIONS IN THIS BLOCK AFTER BACKSUBSTITUTION HAS BEEN
C     CARRIED UP TO BUT NOT INCLUDING EQUATION IPIVOT(LAST).
C     MEANS THAT X(J) CONTAINS THE RIGHT SIDE OF EQUATION IPI-
C     VOT(J) AS MODIFIED DURING ELIMINATION, J=1, ..., LAST, WHILE
C     FOR J .GT. LAST, X(J) IS ALREADY A COMPONENT OF THE SOLUT-
C     ION VECTOR.
C     X(1), ..., X(NCOL) CONTAINS, ON OUTPUT, THE COMPONENTS OF THE SOLUT-
C     ION CORRESPONDING TO THE PRESENT BLOCK.
C
        INTEGER IPIVOT(NROW), LAST, IP, J, K, KP1
        REAL W(NROW, NCOL), X(NCOL), SUM
        K = LAST
        IP = IPIVOT(K)
        SUM = 0.
        IF (K .EQ. NCOL)                GO TO 4
        KP1 = K+1
2         DO 3 J=KP1, NCOL
3         SUM = W(IP, J)*X(J) + SUM
4         X(K) = (X(K) - SUM)/W(IP, K)
        IF (K .EQ. 1)                RETURN
        KP1 = K
        K = K-1
        IP = IPIVOT(K)
        SUM = 0.
        GO TO 2
        END
        BLK40700
        BLK40800
        BLK40900
        BLK41000
        BLK41100
        BLK41200
        BLK41300
        BLK41400
        BLK41500
        BLK41600
        BLK41700
        BLK41800
        BLK41900
        BLK42000
        BLK42100
        BLK42200
        BLK42300
        BLK42400
        BLK42500
        BLK42600
        BLK42700
        BLK42800
        BLK42900
        BLK43000
        BLK43100
        BLK43200
        BLK43300
        BLK43400
        BLK43500
        BLK43600
        BLK43700
        BLK43800

        SUBROUTINE DTBLOK ( BLOKS, INTEGS, NBLOKS, IPIVOT, IFLAG,
        *
        *     DETSGN, DETLOG )
C     COMPUTES THE DETERMINANT OF AN ALMOST BLOCK DIAGONAL MATRIX WHOSE
C     PLU FACTORIZATION HAS BEEN OBTAINED PREVIOUSLY IN FCBLOK.
C     *** THE LOGARITHM OF THE DETERMINANT IS COMPUTED INSTEAD OF THE
C     DETERMINANT ITSELF TO AVOID THE DANGER OF OVERFLOW OR UNDERFLOW
C     INHERENT IN THIS CALCULATION.
C
C     PARAMETERS
C     BLOKS, INTEGS, NBLOKS, IPIVOT, IFLAG ARE AS ON RETURN FROM FCBLOK.
C     IN PARTICULAR, IFLAG = (-1)**(NUMBER OF INTERCHANGES DUR-
C     ING FACTORIZATION) IF SUCCESSFUL, OTHERWISE IFLAG = 0.
C     DETSGN ON OUTPUT, CONTAINS THE SIGN OF THE DETERMINANT.
C     DETLOG ON OUTPUT, CONTAINS THE NATURAL LOGARITHM OF THE DETERMI-
C     NANT IF DETERMINANT IS NOT ZERO. OTHERWISE CONTAINS 0.
C
        INTEGER INTEGS(3, NBLOKS), IPIVOT(1), IFLAG, I, INDEXP, IP, K, LAST
        REAL BLOKS(1), DETSGN, DETLOG
C
        DETSGN = IFLAG
        DETLOG = 0.
        IF (IFLAG .EQ. 0)                RETURN
        INDEX = 0
        BLK43900
        BLK44000
        BLK44100
        BLK44200
        BLK44300
        BLK44400
        BLK44500
        BLK44600
        BLK44700
        BLK44800
        BLK44900
        BLK45000
        BLK45100
        BLK45200
        BLK45300
        BLK45400
        BLK45500
        BLK45600
        BLK45700
        BLK45800
        BLK45900
        BLK46000
        BLK46100

```

```

INDEXP = 0
DO 2 I=1,NBLOKS
  NROW = INTEGS(1,I)
  LAST = INTEGS(3,I)
  DO 1 K=1,LAST
    IP = INDEX + NROW*(K-1) + IPIVOT(INDEXP+K)
    DETLOG = DETLOG + ALOG(ABS(BLOKS(IP)))
1    DETSGN = DETSGN*SIGN(1.,BLOKS(IP))
    INDEX = NROW*INTEGS(2,I) + INDEX
2    INDEXP = INDEXP + NROW
                                RETURN
END
BLK46200
BLK46300
BLK46400
BLK46500
BLK46600
BLK46700
BLK46800
BLK46900
BLK47000
BLK47100
BLK47200
BLK47300

C TEST PROGRAM FOR THE SOLVEBLOK PACKAGE SOLVES AN 11TH ORDER LINEAR
C ALMOST BLOCK DIAGONAL SYSTEM.
  DIMENSION BLOKS(61),B(16),X(11), IPIVOT(16),INTEGS(15)
C      NROW NCOL LAST
  DATA INTEGS/ 3, 4, 2
2      , 3, 3, 3
3      , 3, 4, 1
4      , 3, 4, 1
5      , 4, 4, 4/
  DATA BLOKS /1.,2,-1., 2.,-2.,3, -.1,-.2,-.3, -.1,4.,.3
2      ,0.,-4,3., 0.,.4,.5, 0.,-5.,-.5
3      ,.6,.5,3., -.6,4.,.4, -.6,.5,-.4, 5.,-.5,.4
4      ,0.,0.,.3, 0.,0.,-.3, 0.,0.,.3, 0.,0.,7.
5      ,0.,0.,.2,6., 0.,0.,-.2,.1, 0.,0.,-.2,-.1,
5      ,0.,0.,8.,-.1/
  DATA B /1.94,3.04,-.83, 0.,-3.54,2.75, 1.32,2.35,1.96,
2      2*0.,1.52, 2*0.,.78,2.40 /
  NBLOKS = 5
  N = 11
  CALL SLVBLK ( BLOKS, INTEGS, NBLOKS, B, IPIVOT, X, IFLAG )
  ERROR = 0.
  DO 10 I=1,N
    B(I) = FLOAT(12-I)/10. - X(I)
10  ERROR = AMAX1(ERROR,ABS(B(I)))
  WRITE (6,610) IFLAG,ERROR, (IPIVOT(I),I=1,16), (I,X(I),B(I),I=1,N)
610 FORMAT (28H CHECKOUT SOLVEBLOK ROUTINES/
*      8H IFLAG =,I2,17H, MAXIMUM ERROR =,E9.4//
*      9H IPIVOT =,5(2X,3I2),I2//
*      27H I X(I) ERROR(I)/(I3,F11.5,E13.5))
  CALL DTBLOK ( BLOKS, INTEGS, NBLOKS, IPIVOT, IFLAG,
*      DETSGN, DETLOG )
  DET = DETSGN*EXP(DETLOG)
  ERROR = DET - 2.418821899E6
  WRITE (6,611) DET,ERROR
611 FORMAT(/14H DETERMINANT =,E15.8,11H ERROR = ,E11.5)
                                STOP
END
BLK47400
BLK47500
BLK47600
BLK47700
BLK47800
BLK47900
BLK48000
BLK48100
BLK48200
BLK48300
BLK48400
BLK48500
BLK48600
BLK48700
BLK48800
BLK48900
BLK49000
BLK49100
BLK49200
BLK49300
BLK49400
BLK49500
BLK49600
BLK49700
BLK49800
BLK49900
BLK50000
BLK50100
BLK50200
BLK50300
BLK50400
BLK50500
BLK50600
BLK50700
BLK50800
BLK50900
BLK51000

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ALGORITHM 547

Fortran Routines for Discrete Cubic Spline Interpolation and Smoothing [E1], [E3]

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Key Words and Phrases: discrete splines, discrete cubic splines, discrete natural splines, interpolation, smoothing
CR Categories: 5.12, 5.13
Language: Fortran

DESCRIPTION

1. Introduction

Two Fortran subroutines, DCSINT and DCSSMO, are presented here for discrete cubic spline interpolation and smoothing. The theory for discrete cubic spline interpolation is given by Lyche in [5, 6]. The theory for discrete natural cubic spline smoothing is given by the author in [2]. For most applications continuous splines (see [3, 4, 7]) are probably more appropriate than discrete splines. Discrete spline interpolation and smoothing are worth considering for problems involving functions defined on discrete equally spaced points, or when difference quotients are available as data rather than derivatives. Possible areas of application are discrete time series analysis, computer routines for plotting data, and actuarial- and demographic-type data analysis.

For both discrete cubic spline interpolation and smoothing we direct our attention to the problem of approximating a function $g(t)$ defined on an interval $[\tau_1, \tau_n]$. The values $g(\tau_i) = g_i$ are specified for $i = 1, 2, \dots, n$, where $\tau_i < \tau_{i+1}$. For many applications and for deriving the equations in Sections 2 and 3, it is convenient to assume that $g(t)$ is defined on a discrete point set $T_M = \{t_0, t_1, \dots, t_M\}$ where $t_j = \tau_1 + jh$ for some fixed step size $h > 0$, and also that the τ_i 's belong to T_M . Neither the theory nor the routines DCSINT and DCSSMO require that the τ_i 's belong to T_M . The only mandatory restriction on the τ_i 's is $\tau_i < \tau_{i+1}$.

Discrete cubic splines defined on $[\tau_1, \tau_n]$ with nodes (or knots) $\Lambda_n = \{\tau_1, \tau_2, \dots, \tau_n\}$ are functions $S(t)$ having the form $S(t) = S_i(t)$ where

$$S_i(t) = g_i + b_i(t - \tau_i) + c_i(t - \tau_i)^2 + d_i(t - \tau_i)^3 \quad (1)$$

for $t \in [\tau_i, \tau_{i+1}]$ (i.e., $S(t)$ is piecewise cubic). The $S_i(t)$'s satisfy the joining

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conditions

$$S_i(\tau_{i+1}) = S_{i+1}(\tau_{i+1}), \quad (2)$$

$$\frac{1}{2}(\nabla + \Delta)S_i(\tau_{i+1}) = \frac{1}{2}(\nabla + \Delta)S_{i+1}(\tau_{i+1}), \quad (3)$$

$$\nabla\Delta S_i(\tau_{i+1}) = \nabla\Delta S_{i+1}(\tau_{i+1}), \quad (4)$$

where $\Delta f(x) = f(x+h) - f(x)$ and $\nabla f(x) = f(x) - f(x-h)$. Conditions (3) and (4) should be recognized as requiring the matching up of the first central and the second central differences at τ_{i+1} for adjacent cubics.

Subroutine DCSINT constructs the discrete cubic spline function on the interval $[\tau_1, \tau_n]$ which interpolates the data (τ_j, g_j) for $j = 1, 2, \dots, n$ and satisfies one of the three following end conditions:

I. The first central divided difference end conditions

$$\frac{1}{2} \frac{(\nabla + \Delta)S(\tau_1)}{h} = S_1^{(1)} \quad \text{and} \quad \frac{1}{2} \frac{(\nabla + \Delta)S(\tau_n)}{h} = S_n^{(1)}$$

for specified values of $S_1^{(1)}$ and $S_n^{(1)}$.

II. The second central divided difference end conditions

$$\frac{\nabla\Delta S(\tau_1)}{h^2} = S_1^{(2)} \quad \text{and} \quad \frac{\nabla\Delta S(\tau_n)}{h^2} = S_n^{(2)}$$

for specified values of $S_1^{(2)}$ and $S_n^{(2)}$.

III. The periodic end conditions

$$\frac{1}{2}(\nabla + \Delta)S(\tau_1) = \frac{1}{2}(\nabla + \Delta)S(\tau_n) \quad \text{and} \quad \nabla\Delta S(\tau_1) = \nabla\Delta S(\tau_n)$$

where we assume the data are periodic ($g_1 = g_n$).

In formulating discrete natural cubic spline smoothing, we assume $g(t)$ is defined on a discrete point set T_M . Subroutine DCSSMO constructs the function $S(t)$ which smooths through the data (τ_i, g_i) for $i = 1, 2, \dots, n$ in the sense that (see [3, 7, 8, 9])

$$\sigma(f(t)) = \rho \sum_{i=1}^n W_i [f(\tau_i) - g_i]^2 + \sum_{j=1}^{M-1} [\nabla\Delta f(t_j)]^2 \quad (5)$$

is minimized for $f(t) = S(t)$. This $S(t)$ turns out to be a discrete natural cubic spline ($\nabla\Delta S(\tau_1) = \nabla\Delta S(\tau_n) = 0$). In (5) the W_i 's are positive weights specifying the relative importance of the data, and $\rho > 0$ specifies the amount of smoothing in comparison to data fitting. If ρ is small, the smoothness of the approximation is emphasized at the expense of closely fitting the data. If ρ is large, the discrete natural cubic spline more closely fits the data at the expense of smoothness. In particular, as ρ approaches infinity, the smoothing discrete natural spline approaches the interpolating discrete natural spline.

Sections 2 and 3 describe the equations for constructing the discrete cubic splines for interpolating and smoothing, respectively. Those readers mainly interested in seeing how to use subroutines DCSINT and DCSSMO may turn directly to Sections 4 and 5. Section 6 describes how to construct test examples, and Section 7 presents estimates for execution times.

2. Discrete Cubic Spline Interpolation

The discrete cubic spline $S(t) = S_i(t)$ for $t \in [\tau_i, \tau_{i+1}]$ for $i = 1, 2, \dots, n-1$ where $S_i(t)$ is given in (1). A better form for deriving equations is to represent $S_i(t)$ using factorial polynomials as follows:

$$\begin{aligned} S_i(t) = & g_i + \bar{b}_i(t - \tau_i) + c_i(t - \tau_i)(t - \tau_i - h) \\ & + d_i(t - \tau_i + h)(t - \tau_i)(t - \tau_i - h). \end{aligned} \quad (6)$$

The g_i , c_i , and d_i in (6) are the same as those found in (1), but

$$b_i = \bar{b}_i - c_i h - d_i h^2. \quad (7)$$

The joining conditions (2), (3), and (4) together with $S(\tau_n) = g_n$ give the equations

$$\gamma_i c_i + (\eta_i + \eta_{i+1}) c_{i+1} + \gamma_{i+1} c_{i+2} = 3[g_{i+1}^{(1)} - g_i^{(1)}] \quad (8)$$

for $i = 1, 2, \dots, n-3$, where $H_i = \tau_{i+1} - \tau_i$,

$$\gamma_i = H_i - \frac{h^2}{H_i}, \quad (9)$$

$$\eta_i = 2H_i + \frac{h^2}{H_i}, \quad (10)$$

$$g_i^{(1)} = g[\tau_{i+1}, \tau_i] = \frac{(g_{i+1} - g_i)}{H_i}. \quad (11)$$

The remaining linear equations for the c_i 's come from the end conditions. The complete system of linear equations denoted by

$$Ac = g \quad (12)$$

is now described for the three types of end conditions, I, II, and III. These systems are modified forms of the systems found in Lyche [6].

I. The first central divided difference end conditions have

$$A = \begin{bmatrix} \eta_1 & \gamma_1 & 0 & \dots & 0 \\ \gamma_1 & (\eta_1 + \eta_2) & \gamma_2 & & \\ 0 & \gamma_2 & (\eta_2 + \eta_3) & \dots & \\ \vdots & & \ddots & & 0 \\ \vdots & & & (\eta_{n-2} + \eta_{n-1}) & \gamma_{n-1} \\ 0 & \dots & \gamma_{n-2} & \gamma_{n-1} & \eta_{n-1} \end{bmatrix} \quad (13a)$$

$$c = (c_1, c_2, \dots, c_n)^T, \quad (13b)$$

$$g = 3(g_1^{(1)} - S_1^{(1)}, g_2^{(1)} - g_1^{(1)}, \dots, g_{n-1}^{(1)} - g_{n-2}^{(1)}, S_n^{(1)} - g_{n-1}^{(1)})^T. \quad (13c)$$

The c_n in (13b) is $\nabla \Delta S(\tau_n)/2h^2$ and is introduced to preserve symmetry.

II. The second central divided difference end conditions have

$$A = \begin{bmatrix} (\eta_1 + \eta_2) & \gamma_2 & 0 & \dots & 0 \\ \gamma_2 & (\eta_2 + \eta_3) & \gamma_3 & & \vdots \\ 0 & \gamma_3 & & \ddots & \\ \vdots & & \ddots & & 0 \\ \vdots & & & \gamma_{n-2} & \\ 0 & \dots & 0 & \gamma_{n-2} & (\eta_{n-2} + \eta_{n-1}) \end{bmatrix} \quad (14a)$$

$$c = (c_2, c_3, \dots, c_{n-1})^T, \quad (14b)$$

$$g = (g_2^{(1)} - g_1^{(1)} - \frac{1}{6}\gamma_1 S_1^{(2)}, g_3^{(1)} - g_2^{(1)}, \dots, g_{n-1}^{(1)} - g_{n-2}^{(1)} - \frac{1}{6}\gamma_{n-1} S_{n-1}^{(2)})^T. \quad (14c)$$

In [2] it is shown that the natural end conditions $S_1^{(2)} = S_n^{(2)} = 0$ produce the function which interpolates the given data and minimize $\sum_{i=1}^{n-1} [\nabla \Delta w(t_i)]^2$.

III. The periodic end conditions have

$$A = \begin{bmatrix} (\eta_{n-1} + \eta_1) & \gamma_1 & 0 & \dots & 0 & \gamma_{n-1} \\ \gamma_1 & (\eta_1 + \eta_2) & \gamma_2 & & & 0 \\ 0 & \gamma_2 & & & & \vdots \\ \vdots & & \ddots & & & 0 \\ \vdots & & & \ddots & & \gamma_{n-2} \\ \gamma_{n-1} & 0 & \dots & 0 & \gamma_{n-2} & (\eta_{n-2} + \eta_{n-1}) \end{bmatrix} \quad (15a)$$

$$c = (c_1, c_2, \dots, c_{n-1})^T, \quad (15b)$$

$$g = 3(g_1^{(1)} - g_{n-1}^{(1)}, g_2^{(1)} - g_1^{(1)}, \dots, g_{n-1}^{(1)} - g_{n-2}^{(1)}). \quad (15c)$$

The matrix of the system is tridiagonal except for the γ_{n-1} in the remote corners. This type of linear system is solved in DCSINT using a method described by Björck and Golub in [1], which uses a rank one modification to make (15a) tridiagonal. The number of multiplications and divisions needed to solve system (15) is about $9n$.

The linear systems arising for end conditions I, II, and III are all positive definite and symmetric with nonnegative coefficient matrices A . In all three cases the matrix is strictly diagonally dominant. The infinity condition number for the matrix A in (14a) and (15a) is bounded by (see [2])

$$\text{cond}_\infty(A) \leq \|A\|_\infty \|A^{-1}\| \leq 3 \frac{\max_i(H_{i+1} + H_i)}{\min_i(H_{i+1} + H_i)}. \quad (16a)$$

For (13a) the bound is

$$\text{cond}_\infty(A) \leq 3 \frac{\max_i(H_{i+1} + H_i)}{\min_i(H_i)}. \quad (16b)$$

Hence the A matrices are normally very well conditioned.

Once the c_i 's are known, the b_i and d_i in (1) are given by

$$d_i = \frac{1}{3H_i} [c_{i+1} - c_i], \quad (17)$$

$$b_i = g_i^{(1)} - \frac{H_i}{3} (2c_i + c_{i+1}) \quad (18)$$

for $i = 1, 2, \dots, n-1$ (for the periodic case $c_n = c_1$).

In [5] Lyche gives error bounds for discrete cubic spline interpolation using (I) the first central divided difference end conditions. These bounds involve differences for $g(t)$. He also gives an $O(h^2)$ bound for the distance between the discrete cubic spline and the continuous cubic spline interpolating the same data.

3. Discrete Cubic Spline Smoothing

In [2] the author develops the theory for discrete natural spline smoothing. This type of smoothing was originally studied by Whittaker [8] in a slightly different form (see also [9, pp. 303-316]). The equations for the cubic case can be derived without too much difficulty by developing the cubic discrete analog of Theorem 14.1 given by Greville in [3]. This discrete analog is now stated.

THEOREM 3.1. *Let ρ, h , and $W_i, i = 1, 2, \dots, n$, be positive real numbers. Then there exists a unique function $S(t)$ defined on $T_M = \{t_0, t_1, \dots, t_M\}$ which minimizes*

$$\sigma(f(t)) = \rho \sum_{j=1}^n W_j (f(\tau_j) - g_j)^2 + \sum_{i=1}^{M-1} [\nabla \Delta f(t_i)]^2. \quad (19)$$

This $S(t)$ is the unique discrete natural cubic spline (see (1) or (6)) satisfying

$$S(\tau_i) = g_i - \frac{6h^3}{\rho W_i} [d_i - d_{i-1}] \quad (20)$$

where $d_i = (c_{i+1} - c_i)/(3H_i)$ and $d_0 = d_n = 0$.

The resulting system of linear equations for the c_i 's in (1) or (6) is obtained from (14a-c) and (17) by putting $S_1^{(2)} = S_n^{(2)} = 0$ and replacing the g_i 's used in computing the $g_i^{(1)}$'s by $\tilde{g}_i = S(\tau_i)$, where $S(\tau_i)$ is given by (21).

This linear system is five banded, positive definite, and symmetric (see [2]). The following equations describe this system. For $i = 2, 3, \dots, n-1$ with the

understanding that $c_0 = c_1 = c_n = c_{n+1} = 0$,

$$\begin{aligned} & \beta_{i-1}c_{i-2} + [\gamma_{i-1} - (\beta_{i-1} + \beta_i + \epsilon_{i-1})]c_{i-1} + [(\eta_{i-1} + \eta_i) + (2\beta_i + \epsilon_{i-1} + \epsilon_i)]c_i \\ & + [\gamma_i - (\beta_i + \beta_{i+1} + \epsilon_i)]c_{i+1} + \beta_{i+1}c_{i+2} = 3(g_i^{(1)} - g_{i-1}^{(1)}). \end{aligned} \quad (21)$$

The γ_i , η_i , and $g_i(1)$ are given by (9)-(11),

$$\beta_i = \frac{6h^3}{\rho W_i H_{i-1} H_i} \quad (22)$$

for $i = 1, 2, \dots, n$ with $H_0 = H_n = 1$ ($H_i = \tau_{i+1} - \tau_i$), and

$$\epsilon_i = \frac{(\beta_i H_{i-1} + \beta_{i+1} H_{i+1})}{H_i} \quad (23)$$

for $i = 1, 2, \dots, n-1$.

The discrete cubic spline smoothing the data (τ_i, g_i) for $i = 1, 2, \dots, n$ has the form $S(t) = S_i(t)$ for $t \in [\tau_i, \tau_{i+1}]$ with

$$S_i(t) = \tilde{g}_i + b_i(t - \tau_i) + c_i(t - \tau_i)^2 + d_i(t - \tau_i)^3 \quad (24)$$

where the c_i 's solve (22), (17) gives d_i , (18) gives b_i , and $\tilde{g}_i = S(\tau_i)$ is given in (21).

The five-banded, positive definite, symmetric matrix arising from (22) need not be strictly diagonal dominant, as was the case for interpolation. The condition number of this matrix becomes large as ρ approaches 0. This is to be expected, since the solution to the smoothing problem is not unique when $\rho = 0$. From (23) and (24) we can see that as ρ becomes large, (22) takes the form of the equations for interpolation. Hence (22) is still reasonably conditioned for many useful values of ρ .

4. Examples for DCSINT

We now show how DCSINT is used to find the discrete cubic spline interpolating the data ($h = 0.1$, $n = 6$)

i	τ_i	g_i
1	0.5	1.0
2	0.7	0.5
3	1.0	2.0
4	1.5	2.5
5	2.1	2.0
6	2.5	1.0

for the three end conditions:

I. First central divided difference

$$\frac{1}{2} \frac{(\nabla + \Delta)S(\tau_1)}{h} = -1.0, \quad \frac{1}{2} \frac{(\nabla + \Delta)S(\tau_n)}{h} = 0.0.$$

II. Second central divided difference

$$\frac{\nabla \Delta S(\tau_1)}{h^2} = 0.0, \quad \frac{\nabla \Delta S(\tau_n)}{h^2} = 0.0.$$

(These are the natural end conditions.)

III. Periodic

$$S(\tau_1 - h) = S(\tau_n - h), \quad S(\tau_1) = S(\tau_n), \quad S(\tau_1 + h) = S(\tau_n + h).$$

A possible dimension statement for the calling program is

DIMENSION TNODE (10), G(10), B(10), C(10), D(10)

The nodes τ_i are stored in array TNODE. Note that TNODE (I) must be less than TNODE (I + 1). The data values g_i are stored in array G, and N = 6 and H = 0.1.

I. First Central Divided Difference End Conditions

Put IENT = 1, ENDI = -1.0, and ENDN = 0.0 and call the subroutine as follows:

CALL DCSINT (IENT, H, N, TNODE, G, ENDI, ENDN, B, C, D)

The solution is found in arrays G, B, C, and D.

Solution for I:

I	Interval	G(I)	B(I)	C(I)	D(I)
1	[0.5, 0.7]	1.0	-1.751037	-18.76556	75.10376
2	[0.7, 1.0]	0.5	0.9306393	2.629669	-42.44052
3	[1.0, 1.5]	2.0	4.736782	-11.89979	8.852451
4	[1.5, 2.1]	2.5	-0.4001307	1.378888	-3.501484
5	[2.1, 2.5]	2.0	-2.697525	-4.923788	13.54401

In particular, the discrete cubic spline $S(\tau)$ for $\tau = 1.2$ is given by

$$S(1.2) = 2.0 + 4.736782 \times T - 11.89979 \times T^2 + 8.852451 \times T^3 \quad (25)$$

where $T = (1.2 - 1.0)$

II. Second Central Divided Difference End Conditions

Put IENT = 2, ENDI = 0.0, and ENDN = 0.0, and call DCSINT as before.

Solution for II:

I	G(I)	B(I)	C(I)	D(I)
1	1.0	-4.069061	0.0	39.22655
2	0.5	1.416796	23.53593	-38.63974
3	2.0	4.640036	-11.23984	7.9195391
4	2.5	-0.5627829	0.6394691	-1.817308
5	2.0	-1.798219	-2.631689	2.193076

III. Periodic End Conditions

Put IENT = 3. ENDI and ENDN may contain anything. Call DCSINT as before. (The subroutine assumes $G(1) = G(N)$, so the content of $G(N)$ is never used.)

Solution for III:

I	G(I)	B(I)	C(I)	D(I)
1	1.0	-3.801780	-2.156043	43.32474
2	0.5	1.357985	23.83881	-38.99586
3	2.0	4.663832	-11.25748	7.8596271
4	2.5	-0.6051248	0.5319628	-1.520514
5	2.0	-1.624539	-2.204966	0.04076882

5. An Example for DCSSMO

We now show how DCSSMO is used to find the discrete cubic natural spline which smoothes through the data used in Section 4. No end conditions are needed. A possible dimension statement for the calling program is

DIMENSION TNODE(10), G(10), GS(10), B(10), C(10), D(10), WGS(10)

The subroutine was run for $\rho = \text{RHO} = 0.01$ and for $\rho = \text{RHO} = 1.0$. The weights W_i , stored in array WGS, were taken to be $W_1 = W_5 = W_6 = 2.0$, $W_2 = 2.5$, $W_3 = 1.0$, and $W_4 = 1.5$. The subsolution is called as follows:

CALL DCSSMO (H, N, TNODE, G, WGS, RHO, GS, B, C, D)

The two solutions are as follows.

$\rho = 0.01$:

I	GS(I)	B(I)	C(I)	D(I)
1	0.7307289	1.536621	0.0	0.8975704
2	1.045234	1.667059	0.5385422	-1.374241
3	1.556716	1.611752	-0.6982756	-0.6354344
4	2.108594	0.4271141	-1.651427	0.3430809
5	1.844454	-1.189260	-1.033881	0.8615685

$\rho = 1.0$:

I	GS(I)	B(I)	C(I)	D(I)
1	0.9166568	-2.545056	0.0	27.78108
2	0.6298941	1.329906	16.66866	-26.34149
3	1.817824	3.915267	-7.038706	4.021049
4	2.518413	-0.06161821	-1.007132	-0.5822141
5	1.993115	-1.921916	-2.055118	1.712599

Theoretically, the discrete cubic spline smoothing the given data will approach, in the limit as ρ goes to infinity, the discrete natural cubic spline interpolating these data. From the graphs given in Figure 1 we can see that the discrete smoothing spline for $\rho = 1.0$ more closely follows the discrete natural interpolating spline than does the discrete smoothing spline for $\rho = 0.01$.

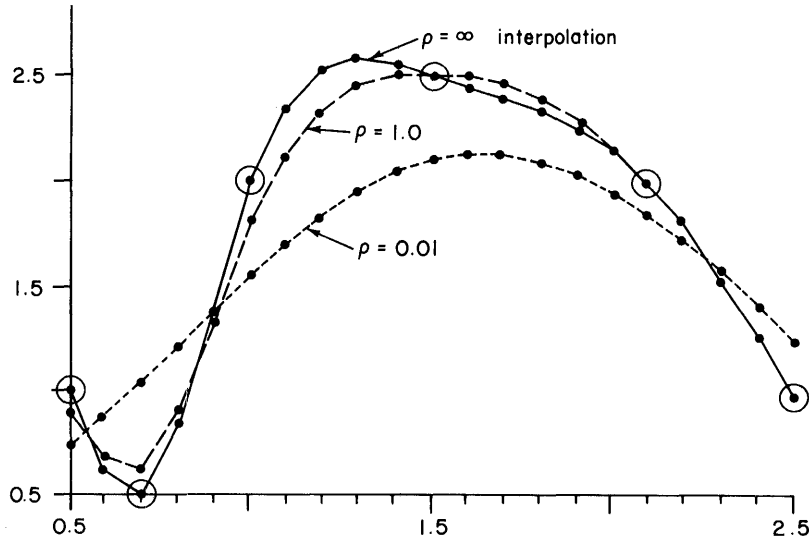


Fig. 1

6. Test Data

Any values may be used for test data $g_i, i = 1, 2, \dots, n$, when interpolating. The correctness of the interpolating solution provided by DCSINT can be checked by verifying that the end conditions are satisfied and that

$$S_i(t) = S_{i+1}(t) \tag{26}$$

for $t = \tau_{i+1} - h, \tau_{i+1}$, and $\tau_{i+1} + h$.

Test data for smoothing (DCSSMO) can be constructed from the discrete natural cubic spline represented in the form

$$S(t) = \alpha_0 + \alpha_1(t - \tau_1) + \sum_{i=1}^n \beta_i(t - \tau_i - h)(t - \tau_i)(t - \tau_i + h)\phi_i(t) \tag{27}$$

where

$$\phi_i(t) = \begin{cases} 0 & \text{for } t < \tau_i \\ 1 & \text{for } t \geq \tau_i \end{cases}$$

and the β_i 's satisfy

$$\sum_{i=1}^n \beta_i = 0, \quad \sum_{i=1}^n \beta_i \tau_i = 0. \tag{28}$$

The data values g_i which produce the $S(t)$ in (28) as the solution to the smoothing problem are given by

$$g_i = S(\tau_i) + 3! \frac{\beta_i}{(\rho W_i)} \tag{29}$$

for $i = 1, 2, \dots, n$. This follows from Theorem 3.1, eq. (20).

7. Timing Estimates

DCSINT and DCSSMO have been run on the IBM 370/168 computer at the UNI-COLL Corporation in Philadelphia. The subroutines have been compiled and run using the WATFIV compiler and the Fortran G and H compilers. Timing estimates were made for the codes compiled by the G and H compilers. These are now listed.

DCSINT

For IENT = 1 (first central divided difference end conditions),

$$T_1 \approx 5.3N \times 10^{-5} \text{ seconds.}$$

For IENT = 2 (second central divided difference end conditions),

$$T_2 \approx 5.7N \times 10^{-5} \text{ seconds.}$$

For IENT = 3 (periodic end conditions),

$$T_3 \approx 7.3N \times 10^{-5} \text{ seconds.}$$

The above timing estimates are about right for the G compiler but are approximately 20 or 25 percent too large for the H compiler using the two-mode optimization parameter.

DCSSMO

$$T_s \approx 11N \times 10^{-5} \text{ seconds.}$$

Again this estimate is right for the G compiler but about 20 or 25 percent too large for the H compiler.

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ALGORITHM

SUBROUTINE DCSSMO(H, N, TNODE, G, WGS, RHO, GSMO, B, C, D)	DCS	10
C	DCS	20
C THIS SUBROUTINE COMPUTES THE DISCRETE NATURAL CUBIC	DCS	30
C SPLINE DEFINED ON THE INTERVAL (TNODE(1),TNODE(N)) WHICH	DCS	40
C SMOOTHS THROUGH THE DATA (TNODE(I),G(I)),I=1,2,...,N.	DCS	50
C N MUST BE 2 OR GREATER. THE NODES MUST SATISFY TNODE(I)	DCS	60
C .LT.TNODE(I+1). THE SOLUTION S(T) FOR T IN THE INTERVAL	DCS	70
C (TNODE(I),TNODE(I+1)) IS GIVEN BY	DCS	80

```

C
C      S(T)=GSMO(I)+B(I)*(T-TNODE(I))+
C          C(I)*(T-TNODE(I))**2+D(I)*(T-TNODE(I))**3
C
C      DIMENSION TNODE(N), G(N), WGS(N), GSMO(N), B(N), C(N), D(N)
C
C INPUT  PARAMETERS(NONE OF THE INPUT PARAMETERS ARE CHANGED
C        BY THIS SUBROUTINE)
C
C H      - THE STEP SIZE USED FOR THE DISCRETE CUBIC SPLINE
C N      - NUMBER OF NODES (TNODE) AND DATA VALUES(G)
C TNODE  - REAL ARRAY CONTAINING THE NODES (TNODE(I).LT.
C          TNODE(I+1)).
C G      - REAL ARRAY CONTAINING THE DATA VALUES.
C WGS    - REAL ARRAY CONTAINING THE WEIGHTS WGS(I)
C          CORRESPONDING TO THE DATA (TNODE(I),G(I)).
C RHO    - SIMPLE REAL VARIABLE CONTAINING THE POSITIVE
C          PARAMETER FOR VARYING THE SMOOTHNESS OF THE FIT.
C          IF RHO IS SMALL SMOOTHNESS IS EMPHASIZED.
C          IF RHO IS LARGE DATA FITTING IS EMPHASIZED.
C
C OUTPUT PARAMETERS
C
C GSMO  - REAL ARRAY CONTAINING THE SMOOTHED VALUES OF
C          THE DATA G(I),I=1,2,...,N.
C B      - REAL ARRAY CONTAINING THE COEFFICIENTS B(I) FOR
C          THE TERMS (T-TNODE(I)).
C C      - REAL ARRAY CONTAINING THE COEFFICIENTS C(I) FOR
C          THE TERMS (T-TNODE(I))**2.
C D      - REAL ARRAY CONTAINING THE COEFFICIENTS D(I) FOR
C          THE TERMS (T-TNODE(I))**3.
C
C      IF (N.EQ.2) GO TO 180
C      N1 = N - 1
C      N2 = N1 - 1
C      N3 = N2 - 1
C THE RIGHT HAND SIDE OF THE LINEAR SYSTEM FOR THE
C C(I)'S WILL NOW BE CONSTRUCTED.
C      DO 10 I=1,N
C          C(I) = G(I)
C 10 CONTINUE
C      DO 20 I=1,N1
C          C(I) = (C(I+1)-C(I))/(TNODE(I+1)-TNODE(I))
C 20 CONTINUE
C      DO 30 I=1,N2
C          C(I) = 3.0*(C(I+1)-C(I))
C 30 CONTINUE
C THE RIGHT HAND SIDE IS NOW IN ARRAY C.
C
C THE P.D. 5 BANDED SYMMETRIC MATRIX WILL NOW BE CONSTRUCTED.
C THE THREE NEEDED DIAGONALS WILL BE STORED IN ARRAYS
C GSMO,B,D.
C      H2 = H*H
C      H3 = H2*H
C      R6 = 6.0*H3/RHO
C      HI3 = TNODE(2) - TNODE(1)
C      HI4 = TNODE(3) - TNODE(2)
C      ETA3 = HI3 + HI3 + H2/HI3
C      BETA3 = R6/(WGS(1)*HI3)
C      BETA4 = R6/(WGS(2)*HI3*HI4)
C      EPS3 = (BETA3+BETA4*HI4)/HI3
C      H2DHI = H2/HI4
C      ETA4 = HI4 + HI4 + H2DHI
C      IF (N.EQ.3) GO TO 60
C      HI5 = TNODE(4) - TNODE(3)
C      BETA5 = R6/(WGS(3)*HI4*HI5)
C      EPS4 = (BETA4*HI3+BETA5*HI5)/HI4
C      GSMO(1) = ETA3 + ETA4 + BETA4 + BETA4 + EPS3 + EPS4
C      P = H2DHI + BETA4 + BETA5 + EPS4
C      B(1) = HI4 - P
C      IF (N.EQ.4) GO TO 50
C      DO 40 I=2,N3
C          HI3 = HI4
C          HI4 = HI5
C          HI5 = TNODE(I+3) - TNODE(I+2)
C          ETA3 = ETA4

```

DCS 90
DCS 100
DCS 110
DCS 120
DCS 130
DCS 140
DCS 150
DCS 160
DCS 170
DCS 180
DCS 190
DCS 200
DCS 210
DCS 220
DCS 230
DCS 240
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DCS 410
DCS 420
DCS 430
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DCS 450
DCS 460
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DCS 490
DCS 500
DCS 510
DCS 520
DCS 530
DCS 540
DCS 550
DCS 560
DCS 570
DCS 580
DCS 590
DCS 600
DCS 610
DCS 620
DCS 630
DCS 640
DCS 650
DCS 660
DCS 670
DCS 680
DCS 690
DCS 700
DCS 710
DCS 720
DCS 730
DCS 740
DCS 750
DCS 760
DCS 770
DCS 780
DCS 790
DCS 800
DCS 810
DCS 820
DCS 830
DCS 840

H2DHI = H2/HI4	DCS 850
ETA4 = HI4 + HI4 + H2DHI	DCS 860
BETA3 = BETA4	DCS 870
BETA4 = BETA5	DCS 880
BETA5 = R6/(WGS(I+2)*HI4*HI5)	DCS 890
EPS3 = EPS4	DCS 900
EPS4 = (BETA4*HI3+BETA5*HI5)/HI4	DCS 910
D(I-1) = BETA4	DCS 920
P = H2DHI + BETA4 + BETA5 + EPS4	DCS 930
B(I) = HI4 - P	DCS 940
GSMO(I) = ETA3 + ETA4 + BETA4 + BETA4 + EPS3 + EPS4	DCS 950
40 CONTINUE	DCS 960
50 HI3 = HI4	DCS 970
HI4 = HI5	DCS 980
ETA3 = ETA4	DCS 990
ETA4 = HI4 + HI4 + H2/HI4	DCS 1000
BETA4 = BETA5	DCS 1010
EPS3 = EPS4	DCS 1020
60 BETA5 = R6/(WGS(N)*HI4)	DCS 1030
EPS4 = (BETA4*HI3+BETA5)/HI4	DCS 1040
GSMO(N2) = ETA3 + ETA4 + BETA4 + BETA4 + EPS3 + EPS4	DCS 1050
C THE P.D. 5 BANDED SYMMETRIC MATRIX IS COMPLETE.	DCS 1060
C THE SYSTEM OF LINEAR EQUATION WILL NOW BE SOLVED FOR THE	DCS 1070
C C(I)'S.	DCS 1080
IF (N.GT.3) GO TO 70	DCS 1090
C(1) = C(1)/GSMO(1)	DCS 1100
GO TO 150	DCS 1110
70 IF (N.GT.4) GO TO 80	DCS 1120
C(1) = (C(1)*GSMO(2)-C(2)*B(1))/(GSMO(1)*GSMO(2)-B(1)**2)	DCS 1130
C(2) = (C(2)-C(1)*B(1))/GSMO(2)	DCS 1140
GO TO 150	DCS 1150
C THIS SOLVE THE 5 BANDED SYSTEM WHEN K=N-2.GT.3.	DCS 1160
80 K = N2	DCS 1170
K1 = K - 1	DCS 1180
K2 = K1 - 1	DCS 1190
K3 = K2 - 1	DCS 1200
C THE 5 BANDED MATRIX WILL NOW BE FACTORED.	DCS 1210
B(1) = B(1)/GSMO(1)	DCS 1220
D(1) = D(1)/GSMO(1)	DCS 1230
P = GSMO(1)*B(1)	DCS 1240
GSMO(2) = GSMO(2) - P*B(1)	DCS 1250
B(2) = (B(2)-P*D(1))/GSMO(2)	DCS 1260
IF (K.EQ.3) GO TO 110	DCS 1270
D(2) = D(2)/GSMO(2)	DCS 1280
IF (K.EQ.4) GO TO 100	DCS 1290
DO 90 I=3,K2	DCS 1300
I1 = I - 1	DCS 1310
I2 = I1 - 1	DCS 1320
P = GSMO(I1)*B(I1)	DCS 1330
GSMO(I) = GSMO(I) - GSMO(I2)*(D(I2)**2) - P*B(I1)	DCS 1340
B(I) = (B(I)-P*D(I1))/GSMO(I)	DCS 1350
D(I) = D(I)/GSMO(I)	DCS 1360
90 CONTINUE	DCS 1370
100 P = GSMO(K2)*B(K2)	DCS 1380
GSMO(K1) = GSMO(K1) - GSMO(K3)*(D(K3)**2) - P*B(K2)	DCS 1390
B(K1) = (B(K1)-P*D(K2))/GSMO(K1)	DCS 1400
110 GSMO(K) = GSMO(K) - GSMO(K2)*(D(K2)**2) - GSMO(K1)*(B(K1)**2)	DCS 1410
C FACTORIZATION COMPLETE.	DCS 1420
C CARRY OUT FORWARD AND BACKWARD SUBSTITUTION.	DCS 1430
C(2) = C(2) - B(1)*C(1)	DCS 1440
DO 120 I=3,K	DCS 1450
I1 = I - 1	DCS 1460
I2 = I - 2	DCS 1470
C(I) = C(I) - B(I1)*C(I1) - D(I2)*C(I2)	DCS 1480
120 CONTINUE	DCS 1490
DO 130 I=1,K	DCS 1500
C(I) = C(I)/GSMO(I)	DCS 1510
130 CONTINUE	DCS 1520
C(K1) = C(K1) - B(K1)*C(K)	DCS 1530
DO 140 I=2,K1	DCS 1540
J = K - I	DCS 1550
C(J) = C(J) - B(J)*C(J+1) - D(J)*C(J+2)	DCS 1560
140 CONTINUE	DCS 1570
C THE 5 BANDED SYSTEM HAS BEEN SOLVED.THE SOLUTION IS IN	DCS 1580
DO 130 I=1,K	DCS 1590
C(I) = C(I)/GSMO(I)	DCS 1510


```

130 CONTINUE
  C(K1) = C(K1) - B(K1)*C(K)
  DO 140 I=2,K1
    J = K - I
    C(J) = C(J) - B(J)*C(J+1) - D(J)*C(J+2)
140 CONTINUE
C THE 5 BANDED SYSTEM HAS BEEN SOLVED.THE SOLUTION IS IN
C ARRAY C. THE COEFFICIENTS GSMO, B, C, AND D WILL NOW BE
C SET UP.
150 C(N) = 0.0
  D(N) = 0.0
  C(N1) = C(N2)
  HK1 = TNODE(N) - TNODE(N1)
  D(N1) = -C(N1)/(3.0*HK1)
  GSMO(N) = G(N) + R6*D(N1)/WGS(N)
  IF (N.EQ.3) GO TO 170
  DO 160 I=2,N2
    K = N - I
    K1 = K + 1
    HK2 = HK1
    HK1 = TNODE(K1) - TNODE(K)
    C(K) = C(K-1)
    D(K) = (C(K1)-C(K))/(3.0*HK1)
    GSMO(K1) = G(K1) - R6*(D(K1)-D(K))/WGS(K1)
    B(K1) = (GSMO(K1+1)-GSMO(K1))/HK2 - HK2*(C(K1)+C(K1)+C(K1+1))/
      * 3.0
160 CONTINUE
170 C(1) = 0.0
  HK2 = HK1
  HK1 = TNODE(2) - TNODE(1)
  D(1) = (C(2)-C(1))/(3.0*HK1)
  GSMO(2) = G(2) - R6*(D(2)-D(1))/WGS(2)
  GSMO(1) = G(1) - R6*D(1)/WGS(1)
  B(2) = (GSMO(3)-GSMO(2))/HK2 - HK2*(C(2)+C(2)+C(3))/3.0
  B(1) = (GSMO(2)-GSMO(1))/HK1 - HK1*(C(1)+C(1)+C(2))/3.0
C THE DISCRETE CUBIC SMOOTHING SPLINE IS NOW COMPLETE.
  RETURN
C THE TRIVIAL CASE WHEN N=2 IS HANDLED HERE.
180 GSMO(1) = G(1)
  GSMO(2) = G(2)
  B(1) = (G(2)-G(1))/(TNODE(2)-TNODE(1))
  C(1) = 0.0
  D(1) = 0.0
  RETURN
  END

      SUBROUTINE DCSINT(IENT, H, N, TNODE, G, ENDL, ENDN, B, C, D)
C
C THIS SUBROUTINE COMPUTES THE DISCRETE CUBIC SPLINE
C DEFINED ON THE INTERVAL (TNODE(1),TNODE(N)),WHICH INTER-
C POLATES THE DATA (TNODE(I),G(I)),I=1,2,...,N. WE REQUIRE
C THAT TNODE(1).LT.TNODE(I+1). ENDL AND ENDN CONTAIN THE
C VALUES OF THE END CONDITIONS BEING USED.
C
C IF IENT=1,THE FIRST CENTRAL DIVIDED DIFFERENCE END
C CONDITIONS ARE BEING USED.
C
C IF IENT=2,THE SECOND CENTRAL DIVIDED DIFFERENCE END
C CONDITIONS ARE BEING USED.
C
C IF IENT=3,THE PERODIC END CONDITIONS ARE BEING USED.
C FOR THIS CASE THE CONTENTS OF G(N), ENDL,AND ENDN ARE
C IGNORED.
C
C FOR ALL THREE END CONDITIONS N MUST BE GREATER THAN OR
C EQUAL TO 2.
C
C THE DISCRETE CUBIC SPLINE IS REPRESENTED BY PIECEWISE
C CUBIC POLYNOMIALS. FOR T IN THE INTERNAL (TNODE(I),TNODE
C (I+1)) THE CUBIC SPLINE IS
C
C      S(T)=G(I)+B(I)*(T-TNODE(I))
C              +C(I)*(T-TNODE(I))**2
C              +D(I)*(T-TNODE(I))**3

```

DCS 1520
DCS 1530
DCS 1540
DCS 1550
DCS 1560
DCS 1570
DCS 1580
DCS 1590
DCS 1600
DCS 1610
DCS 1620
DCS 1630
DCS 1640
DCS 1650
DCS 1660
DCS 1670
DCS 1680
DCS 1690
DCS 1700
DCS 1710
DCS 1720
DCS 1730
DCS 1740
DCS 1750
DCS 1760
DCS 1770
DCS 1780
DCS 1790
DCS 1800
DCS 1810
DCS 1820
DCS 1830
DCS 1840
DCS 1850
DCS 1860
DCS 1870
DCS 1880
DCS 1890
DCS 1900
DCS 1910
DCS 1920
DCS 1930
DCS 1940
DCS 1950
DCS 1960

DCS 10
DCS 20
DCS 30
DCS 40
DCS 50
DCS 60
DCS 70
DCS 80
DCS 90
DCS 100
DCS 110
DCS 120
DCS 130
DCS 140
DCS 150
DCS 160
DCS 170
DCS 180
DCS 190
DCS 200
DCS 210
DCS 220
DCS 230
DCS 240
DCS 250
DCS 260
DCS 270
DCS 280
DCS 290

DIMENSION TNODE(N), G(N), B(N), C(N), D(N)	DCS	300
C	DCS	310
C INPUT PARAMETERS (NONE OF THESE PARAMETERS	DCS	320
C ARE CHANGED BY THIS SUBROUTINE.)	DCS	330
C	DCS	340
C IENT - SPECIFIES END CONDITIONS WHICH ARE IN EFFECT.	DCS	350
C H - THE STEP SIZE USED FOR THE DISCRETE CUBIC SPLINE.	DCS	360
C N - NUMBER OF NODES (TNODE) AND DATA VALUES (G).	DCS	370
C (N.GE.2)	DCS	380
C TNODE - REAL ARRAY CONTAINING THE NODES (TNODE(I).LT.	DCS	390
C TNODE(I+1)).	DCS	400
C G - REAL ARRAY CONTAINING THE INTERPOLATING DATA.	DCS	410
C END1 - END CONDITION VALUE AT TNODE(1).	DCS	420
C ENDN - END CONDITION VALUE AT TNODE(N).	DCS	430
C	DCS	440
C OUTPUT PARAMETERS	DCS	450
C	DCS	460
C B - REAL ARRAY CONTAINING COEFFICIENTS OF	DCS	470
C (T-TNODE(I)), I=1,2,...,N-1.	DCS	480
C C - REAL ARRAY CONTAINING COEFFICIENTS OF	DCS	490
C (T-TNODE(I))**2, I=1,2,...,N-1.	DCS	500
C D - REAL ARRAY CONTAINING COEFFICIENTS OF	DCS	510
C (T-TNODE(I))**3, I=1,2,...,N-1.	DCS	520
C	DCS	530
C SPECIAL CASES ARE ACCOUNTED FOR HERE.	DCS	540
IF (N.EQ.2 .AND. IENT.EQ.3) GO TO 220	DCS	550
H2 = H*H	DCS	560
N1 = N - 1	DCS	570
IF (N.EQ.2 .AND. IENT.EQ.2) GO TO 180	DCS	580
N2 = N1 - 1	DCS	590
C THE SYMMETRIC TRIDIAGONAL(OR NEAR TRIDIAGONAL) LINEAR	DCS	600
C SYSTEM WILL NOW BE SET UP FOR THE APPROPRIATE END	DCS	610
C CONDITIONS	DCS	620
HI = TNODE(2) - TNODE(1)	DCS	630
H2DHI = H2/HI	DCS	640
ETA2 = HI + HI + H2DHI	DCS	650
GO TO (10, 40, 60), IENT	DCS	660
C IENT=1 - FIRST CENTRAL DIVIDED DIFFERENCE END CONDITONS	DCS	670
C SET UP.	DCS	680
10 B(1) = ETA2	DCS	690
D(1) = HI - H2DHI	DCS	700
G2 = (G(2)-G(1))/HI	DCS	710
C(1) = 3.0*(G2-END1)	DCS	720
IF (N.EQ.2) GO TO 30	DCS	730
DO 20 I=2,N1	DCS	740
ETA1 = ETA2	DCS	750
G1 = G2	DCS	760
HI = TNODE(I+1) - TNODE(I)	DCS	770
H2DHI = H2/HI	DCS	780
ETA2 = HI + HI + H2DHI	DCS	790
B(I) = ETA1 + ETA2	DCS	800
D(I) = HI - H2DHI	DCS	810
G2 = (G(I+1)-G(I))/HI	DCS	820
C(I) = 3.0*(G2-G1)	DCS	830
20 CONTINUE	DCS	840
30 B(N) = ETA2	DCS	850
C(N) = 3.0*(ENDN-G2)	DCS	860
L = N	DCS	870
C SET UP FOR (1) FIRST CENTRAL DIVIDED DIFFERENCE END	DCS	880
C CONDITING COMPLETE.THE LINEAR EQUATIONS WILL BE NOW	DCS	890
C SOLVED.	DCS	900
GO TO 110	DCS	910
C IENT=2 - SECOND CENTRAL DIVIDED DIFFERENCE END CONDITIONS	DCS	920
C SET UP.	DCS	930
40 GAMMA = HI - H2DHI	DCS	940
G2 = (G(2)-G(1))/HI	DCS	950
DO 50 I=1,N2	DCS	960
ETA1 = ETA2	DCS	970
G1 = G2	DCS	980
HI = TNODE(I+2) - TNODE(I+1)	DCS	990
H2DHI = H2/HI	DCS	1000
ETA2 = HI + HI + H2DHI	DCS	1010
B(I) = ETA1 + ETA2	DCS	1020
D(I) = HI - H2DHI	DCS	1030
G2 = (G(I+2)-G(I+1))/HI	DCS	1040
C(I) = 3.0*(G2-G1)	DCS	1050

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50 CONTINUE
C(1) = C(1) - GAMMA*END1/2.0
HI = TNODE(N) - TNODE(N1)
GAMMA = HI - H2/HI
C(N2) = C(N2) - GAMMA*ENDN/2.0
C STEP UP FOR (2) SECOND CENTRAL DIVIDED DIFFERENCE
C END CONDITIONS COMPLETE. THE LINEAR EQUATIONS WILL NOW
C BE SOLVED.
IF (N2.EQ.1) GO TO 100
L = N2
GO TO 110
C IENT=3 - PERIODIC END CONDITIONS SET UP.
60 B(1) = ETA2
D(1) = HI - H2DHI
DO 70 I=2,N1
ETA1 = ETA2
HI = TNODE(I+1) - TNODE(I)
H2DHI = H2/HI
ETA2 = HI + HI + H2DHI
B(I) = ETA1 + ETA2
D(I) = HI - H2DHI
C(I) = 0.0
70 CONTINUE
CT = D(N1)
C(1) = CT
C(N1) = CT
B(1) = B(1) + ETA2 - CT
B(N1) = B(N1) - CT
L = N1
ITRANS = 1
GO TO 120
80 G1 = (G(N1)-G(N2))/(TNODE(N1)-TNODE(N2))
G2 = (G(1)-G(N1))/(TNODE(N)-TNODE(N1))
DEN = (1.0+C(1)+C(N1))
CT = 3.0*(G2-G1)
BS1 = CT*C(N1)
C(N1) = CT
DO 90 I=1,N2
HI = TNODE(I+1) - TNODE(I)
G1 = G2
G2 = (G(I+1)-G(I))/HI
CT = 3.0*(G2-G1)
BS1 = BS1 + CT*C(I)
C(I) = CT
90 CONTINUE
BS1 = BS1/DEN
C(1) = C(1) - BS1
C(N1) = C(N1) - BS1
ITRANS = 0
GO TO 140
C THE SET UP AND MOST OF THE DETAILS FOR SOLVING THE
C LINEAR EQUATION FOR (3) THE PERIODIC END CONDITIONS
C ARE COMPLETED.
C
C THE LINEAR EQUATION ARE SOLVED HERE.
100 C(2) = C(1)/B(1)
GO TO 180
110 ITRANS = 0
120 LI = L - 1
DO 130 I=1,LI
T = D(I)/B(I)
B(I+1) = B(I+1) - T*D(I)
D(I) = T
130 CONTINUE
C THE LINEAR EQUATION SOLVER IS ENTERED AT THIS POINT
C IF THE LU FACTORIZATION HAS ALREADY BEEN DONE.
140 DO 150 I=1,LI
C(I+1) = C(I+1) - D(I)*C(I)
150 CONTINUE
C(L) = C(L)/B(L)
DO 160 I=1,LI
LI = L - I
C(LI) = C(LI)/B(LI) - D(LI)*C(LI+1)
160 CONTINUE
IF (ITRANS.GE.1) GO TO 80
C THE LINEAR SYSTEM HAS BEEN SOLVE FOR THE C-VECTOR

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DCS 1060
DCS 1070
DCS 1080
DCS 1090
DCS 1100
DCS 1110
DCS 1120
DCS 1130
DCS 1140
DCS 1150
DCS 1160
DCS 1170
DCS 1180
DCS 1190
DCS 1200
DCS 1210
DCS 1220
DCS 1230
DCS 1240
DCS 1250
DCS 1260
DCS 1270
DCS 1280
DCS 1290
DCS 1300
DCS 1310
DCS 1320
DCS 1330
DCS 1340
DCS 1350
DCS 1360
DCS 1370
DCS 1380
DCS 1390
DCS 1400
DCS 1410
DCS 1420
DCS 1430
DCS 1440
DCS 1450
DCS 1460
DCS 1470
DCS 1480
DCS 1490
DCS 1500
DCS 1510
DCS 1520
DCS 1530
DCS 1540
DCS 1550
DCS 1560
DCS 1570
DCS 1580
DCS 1590
DCS 1600
DCS 1610
DCS 1620
DCS 1630
DCS 1640
DCS 1650
DCS 1660
DCS 1670
DCS 1680
DCS 1690
DCS 1700
DCS 1710
DCS 1720
DCS 1730
DCS 1740
DCS 1750
DCS 1760
DCS 1770
DCS 1780
DCS 1790
DCS 1800
DCS 1810

```

      IF (IENT.EQ.3) C(N) = C(1)
      IF (IENT.NE.2) GO TO 190
      DO 170 I=1,N2
        LI = N - I
        C(LI) = C(LI-1)
170  CONTINUE
180  C(1) = END1/2.0
      C(N) = ENDN/2.0
190  C1 = C(1)
      C2 = C(2)
      HI = TNODE(2) - TNODE(1)
      IF (N.EQ.2) GO TO 210
      DO 200 I=1,N2
        B(I) = (G(I+1)-G(I))/HI - HI*(C1+C1+C2)/3.0
        D(I) = (C2-C1)/(3.0*HI)
        HI = TNODE(I+2) - TNODE(I+1)
        C1 = C2
        C2 = C(I+2)
200  CONTINUE
210  GN = G(1)
      IF (IENT.NE.3) GN = G(N)
      B(N1) = (GN-G(N1))/HI - HI*(C1+C1+C2)/3.0
      D(N1) = (C2-C1)/(3.0*HI)
C THE INTERPOLATING DISCRETE CUBIC SPLINE HAS BEEN
C CONSTRUCTED.
      RETURN
C THE FOLLOWING HANDLES THE TRIVIAL PERIODIC CASE
C (IENT.EQ.3) WHEN N.EQ.2.
220  B(1) = 0.0
      C(1) = 0.0
      D(1) = 0.0
      RETURN
      END

```

DCS 1820
DCS 1830
DCS 1840
DCS 1850
DCS 1860
DCS 1870
DCS 1880
DCS 1890
DCS 1900
DCS 1910
DCS 1920
DCS 1930
DCS 1940
DCS 1950
DCS 1960
DCS 1970
DCS 1980
DCS 1990
DCS 2000
DCS 2010
DCS 2020
DCS 2030
DCS 2040
DCS 2050
DCS 2060
DCS 2070
DCS 2080
DCS 2090
DCS 2100
DCS 2110
DCS 2120
DCS 2130
DCS 2140

```

C TEST FOR DURIS ALGORITHM, JAN 1979
C DRIVER FOR DCSSMO
C THIS DRIVER USES SUBROUTINE DCSSMO TO COMPUTE THE
C DISCRETE NATURAL CUBIC SPLINE WHICH SMOOTHS THROUGH
C THE DATA (TNODE(I),G(I)),I=1,2,...,N AS SPECIFIED
C BY THE SMOOTHING PARAMETER RHO AND THE WEIGHTS WGS(I),
C I=1,2,...,N.
      DIMENSION TNODE(20), G(20), B(20), C(20), D(20), GSMO(20),
      * WGS(20)
C THE NEXT READ STATEMENT BRINGS IN THE TOTAL NUMBER
C OF DATA POINTS N, THE STEP SIZE H, AND THE PARAMETER RHO
C SPECIFYING THE AMOUNT OF SMOOTHING DESIRED.
      READ (5,99999) N, H, RHO
C THE FOLLOW WRITE STATEMENTS OUTPUT N,H,AND RHO AND
C PREPARES TO OUTPUT DATA AND SOLUTION.
      WRITE (6,99998)
      WRITE (6,99997) N, H, RHO
C THE NEXT DO LOOP READS IN THE DATA VALUES
C (TNODE(I),G(I)), AND THE WEIGHTS WGS(I) FOR
C I=1,2,...,N AND AT THE SAME TIME WRITES THIS DATA
C AS OUTPUT TO THE USER.
      DO 10 I=1,N
        READ (5,99995) TNODE(I), G(I), WGS(I)
        WRITE (6,99996) I, TNODE(I), G(I), WGS(I)
10  CONTINUE
C SUBROUTINE DCSSMO PARAMETERS H, TNODE,G,WGS,
C ARE AS SPECIFIED ABOVE.
      DRHO=RHO
      NN=N
      DO 50 N=2,NN
        DO 40 IRHO=1,3
          RHO=DRHO*100.**(IRHO-2)
          WRITE (6,99998)
          WRITE (6,99997) N, H, RHO
          DO 15 I=1,N
            WRITE (6,99996) I, TNODE(I), G(I), WGS(I)
15  CONTINUE
          CALL DCSSMO(H, N, TNODE, G, WGS, RHO, GSMO, B, C, D)
C OUTPUT FROM DCSSMO IS GSMO(I),B(I),C(I),AND D(I) FOR
C I=1,2,...,N-1. IN PARTICULAR THE DISCRETE CUBIC SPLINE
C ON THE INTERVAL (TNODE(1),TNODE(I+1)) IS DESCRIBED BY

```

DU 00010
DU 00020
DU 00030
DU 00040
DU 00050
DU 00060
DU 00070
DU 00080
DU 00090
DU 00100
DU 00110
DU 00120
DU 00130
DU 00140
DU 00150
DU 00160
DU 00170
DU 00180
DU 00190
DU 00200
DU 00210
DU 00220
DU 00230
DU 00240
DU 00250
DU 00260
DU 00270
DU 00280
DU 00290
DU 00300
DU 00310
DU 00320
DU 00330
DU 00340
DU 00350
DU 00360
DU 00370
DU 00380
DU 00390
DU 00400
DU 00410

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C THE CUBIC POLYNOMIAL DU 00420
C DU 00430
C S(T)=GSMO(I)+B(I)*(T-TNODE(I))+ DU 00440
C C(I)*(T-TNODE(I))**2+D(I)*(T-TNODE(I))**3 DU 00450
C DU 00460
C N1 = N - 1 DU 00470
C THE SOLUTION IS NOW PRINTED OUT BY THE FOLLOWING STATEMENTS DU 00480
WRITE (6,99994) DU 00490
DO 20 I=1,N1 DU 00500
WRITE (6,99993) I, GSMO(I), B(I), C(I), D(I) DU 00510
20 CONTINUE DU 00520
WRITE (6,99992) DU 00530
DO 30 I=1,N1 DU 00540
T=TNODE(I+1)-TNODE(I) DU 00550
SL=GSMO(I) DU 00560
SR=GSMO(I)+(B(I)+(C(I)+D(I)*T)*T)*T DU 00570
FL=B(I)+D(I)*H*H DU 00580
FR=FL+(2.*C(I)+3.*D(I)*T)*T DU 00590
DL=2.*C(I) DU 00600
DR=DL+6.*D(I)*T DU 00610
WRITE (6,99993) I, SL, SR, FL, FR, DL, DR DU 00620
30 CONTINUE DU 00630
40 CONTINUE DU 00640
50 CONTINUE DU 00650
CALL TEST2 DU 00660
STOP DU 00670
99999 FORMAT (I10, 2E10.2) DU 00680
99998 FORMAT (///5H DATA) DU 00690
99997 FORMAT (3H N=, I3, 2X, 2HH=, F10.4, 3X, 4HRHO=, F10.4 //5H I, DU 00700
* 3X, 8HTNODE(I), 9X, 4HG(I), 7X, 6HWGS(I)) DU 00710
99996 FORMAT (I3, 3E16.7) DU 00720
99995 FORMAT (3E10.2) DU 00730
99994 FORMAT (///40H DISCRETE NATURAL CUBIC SMOOTHING SPLINE//4H I, DU 00740
* 2X, 7HGSMO(I), 12X, 4HB(I), 16X, 4HC(I), 16X, 4HD(I)) DU 00750
99993 FORMAT (I3, 6E17.7) DU 00760
99992 FORMAT(3H0 I,14X,6HVALUES,18X,25HFIRST DIVIDED DIFFERENCES, DU 00770
1 9X,26HSECOND DIVIDED DIFFERENCES) DU 00780
END DU 00790

SUBROUTINE TEST2 DU 00800
C DRIVER FOR DCSINT DU 00810
C THIS DRIVER USES SUBROUTINE DCSINT TO COMPUTE THE DU 00820
C DISCRETE CUBIC SPLINE WHICH INTERPOLATES THE DATA DU 00830
C (TNODE(I),G(I)),I=1,2,...,N SUBJECT TO THE END CONDITIONS DU 00840
C SPECIFIED BY IENT,END1,AND ENDN. DU 00850
C DU 00860
C DIMENSION TNODE(20), G(20), B(20), C(20), D(20) DU 00870
C DU 00880
C THE NEXT READ STATEMENT BRINGS IN THE NUMBER OF DATA DU 00890
C POINTS N AND THE DISCRETE CUBIC SPLINE STEP SIZE H. DU 00900
C DU 00910
C READ (5,99999) N, H DU 00920
C WRITE (6,99998) N, H DU 00930
C WRITE (6,99997) DU 00940
C THE NEXT DO LOOP READS IN AND WRITES OUT THE DATA DU 00950
C VALUES (TNODE(I),G(I)),I=1,2,...,N. DU 00960
C DO 10 I=1,N DU 00970
C READ (5,99996) TNODE(I), G(I) DU 00980
C WRITE (6,99995) I, TNODE(I), G(I) DU 00990
10 CONTINUE DU 01000
C THE END CONDITIONS ARE NOW READ IN. IENT SPECIFIES DU 01010
C THE TYPE OF END CONDITION(SEE SUBROUTINE COMMENTS) DU 01020
C AND END1 AND ENDN SPECIFY THE VALUES OF THE END CONDITIONS. DU 01030
C READ (5,99994) IENT, END1, ENDN DU 01040
C SUBROUTINE DCSINT IS NOW CALLED. THE PARAMETERS DU 01050
C H,TNODE,G,END1,AND ENDN ARE AS SPECIFIED ABOVE DU 01060
C NN=N DU 01070
C DO 50 N=2,NN DU 01080
C WRITE (6,99998) N, H DU 01090
C WRITE (6,99997) DU 01100
C DO 15 I=1,N DU 01110
C WRITE (6,99995) I, TNODE(I), G(I) DU 01120
15 CONTINUE DU 01130
C DO 40 IENT=1,3 DU 01140
C DU 01150

```

```

      CALL DCSINT(IENT, H, N, TNODE, G, ENDL, ENDN, B, C, D)
C THE END CONDITIONS IENT, ENDL, ENDN AND THE DISCRETE
C CUBIC SPLINE SOLUTION TO THE INTERPOLATIONPROBLEM
C ARE NOW OUTPUTED BELOW. IN PARTICULAR THE DISCRETE
C CUBIC SPLINE ON THE INTERVAL (TNODE(I), TNODE(I+1))
C IS DESCRIBES BY THE CUBIC POLYNOMIAL
C
      S(T)=G(I)+B(I)*ARG+C(I)*ARG**2
      +D(I)*ARG**3
C
C WHERE ARG=(T-TNODE(I)).
C
      WRITE (6,99993) IENT, ENDL, ENDN
      WRITE (6,99992)
      N1 = N - 1
      DO 20 I=1,N1
        WRITE (6,99991) I, G(I), B(I), C(I), D(I)
20 CONTINUE
      WRITE(6,99990)
      DO 30 I=1,N1
        T=TNODE(I+1)-TNODE(I)
        SL=G(I)
        SR=G(I)+(B(I)+(C(I)+D(I)*T)*T)*T
        FL=B(I)+D(I)*H*H
        FR=FL+(2.*C(I)+3.*D(I)*T)*T
        DL=2.*C(I)
        DR=DL+6.*D(I)*T
        WRITE(6,99991) I, SL, SR, FL, FR, DL, DR
30 CONTINUE
40 CONTINUE
50 CONTINUE
      RETURN
99999 FORMAT (I10, E10.2)
99998 FORMAT (///5H DATA//3H N=, I3, 2X, 2HH=, F10.5)
99997 FORMAT (2H I, 3X, 8HTNODE(I), 10X, 4HG(I))
99996 FORMAT (2E10.2)
99995 FORMAT (I3, 3E16.7)
99994 FORMAT (I10, 2E10.2)
99993 FORMAT (///6H IENT=, I2, 2X, 5HEND1=, E16.7, 2X, 5HEND2=, E16.7)
99992 FORMAT (/2H I, 3X, 4HG(I), 13X, 4HB(I), 13X, 4HC(I), 13X, 4HD(I))
99991 FORMAT (I3, 6E17.7)
99990 FORMAT(3H0 I, 14X, 6HVALUES, 18X, 25HFIRST DIVIDED DIFFERENCES,
      1 9X, 26HSECOND DIVIDED DIFFERENCES)
      END

```

	7	0.1	100.0	DU 01610
0.5		1.0	2.0	DU 01620
0.7		0.5	2.5	DU 01630
1.0		2.0	1.0	DU 01640
1.5		2.5	1.5	DU 01650
2.1		2.0	2.0	DU 01660
2.5		1.0	2.0	DU 01670
3.0		0.1	2.0	DU 01680
	6	0.1		DU 01690
0.5		1.0		DU 01700
0.7		0.5		DU 01710
1.0		2.0		DU 01720
1.5		2.5		DU 01730
2.1		2.0		DU 01740
2.5		1.0		DU 01750
	2	0.001	0.002	DU 01760

ALGORITHM 548

Solution of the Assignment Problem [H]

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Key Words and Phrases: assignment problem, Hungarian algorithm

CR Categories: 5.39, 8.3

Language: Fortran

DESCRIPTION

Problem

The algorithm presented in this paper solves the assignment problem of the following form: given an $n \times n$ cost matrix $(a_{i,j})$, find a permutation (C_i) of integers $1, \dots, n$ that minimizes

$$T = \sum_{i=1}^n a_{i,C_i}.$$

It is supposed, without loss of generality, that the elements of the cost matrix are nonnegative integers.

Algorithm

To give the reader a better understanding of the algorithm proposed below, we define:

- C_j as the row assigned to column j ($j = 1, \dots, n$);
- LC_j as the label of column j ; if $LC_j = 0$, column j is unlabeled ($j = 1, \dots, n$);
- LR_i as the label of row i ; if $LR_i = 0$, row i is unlabeled ($i = 1, \dots, n$);
- T as the assignment cost;
- P_i as the set containing the columns corresponding to the unassigned zero elements of row i of the cost matrix ($i = 1, \dots, n$);
- RH as the set containing the current not-completely-explored rows;
- U as the set containing the unassigned rows;
- $first(s)$ as the first element of set s ;
- $next(s)$ as the element following the last considered element of set s ;
- $last(s)$ as the last element of set s .

Step 1. [Initialization]

Set $C_j = 0$, for $j = 1, \dots, n$;
 set $U = \emptyset$, $N = \{i \mid 1 \leq i \leq n\}$.

Step 2. [Reduction of the initial cost matrix]

Set $S_j = \min_{i \in N} [a_{i,j}]$, for all $j \in N$;

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set $Q_i = \min_{j \in N} [a_{i,j} - S_j]$, for all $i \in N$;
 set $a'_{i,j} = a_{i,j} - S_j - Q_i$, for all $i \in N, j \in N$;
 set $P_i = \{j \in N \mid a'_{i,j} = 0\}$, for all $i \in N$; set $T = \sum_{k \in N} (S_k + Q_k)$.

Step 3. [Choice of the initial solution]

Set $i = 1$.

- a. Set $j = \text{first} \{k \in P_i \mid C_k = 0\}$; if j exists, go to Step 3d. Otherwise, set $j = \text{first} \{P_i\}$.
- b. Set $m = \text{first} \{k \in P_{C_j} \mid C_k = 0\}$; if m exists, go to Step 3c. Otherwise, set $j = \text{next} \{P_i\}$; if j exists, repeat Step 3b; if not, set $U = U \cup \{i\}$, go to Step 3e.
- c. Set $C_m = C_j, P_{C_j} = P_{C_j} \cup \{j\} - \{m\}$.
- d. Set $C_j = i, P_i = P_i - \{j\}$.
- e. If $i < n$, set $i = i + 1$, go to Step 3a.

Step 4. [Search for a new assignment]

If set U is empty, stop (the optimal assignment is given by vector (C_j) , the minimum cost is given by T). Otherwise, set $RH = \emptyset$; set $LR_k = LC_k = 0$, for $k = 1, 2, \dots, n$; set $r = \text{first} \{U\}, LR_r = -1$.

- a. If set P_r is empty, go to Step 4c. Otherwise, if set P_r has at least two elements, set $RH = RH \cup \{r\}$; in any case, set $l = \text{first} \{P_r\}$.
- b. If $LC_l = 0$, set $LC_l = r$, if $C_l = 0$, go to Step 6; if not, set $r = C_l, LR_r = l$, go to Step 4a. Otherwise ($LC_l \neq 0$), if $r \in RH$, go to Step 4d.
- c. If set RH is empty, go to Step 5. Otherwise, set $r = \text{first} \{RH\}$.
- d. Set $l = \text{next} \{P_r\}$, if $l = \text{last} \{P_r\}$, set $RH = RH - \{r\}$; in any case, go to Step 4b.

Step 5. [Reduction of the current cost matrix]

Set $SLR = \{i \in N \mid LR_i \neq 0\}, SUC = \{j \in N \mid LC_j = 0\}$.

Set $H = \min_{i \in SLR, j \in SUC} [a'_{i,j}]$.

For all $i \in SLR, j \in SUC$: set $a'_{i,j} = a'_{i,j} - H$, if $a'_{i,j} = 0$, set $P_i = P_i \cup \{j\}, RH = RH \cup \{i\}$;

For all $i \in N - SLR, j \in N - SUC$: if $j \in P_i$, set $P_i = P_i - \{j\}$; in any case set $a'_{i,j} = a'_{i,j} + H$.

Set $T = T + H, r = \text{first} \{RH\}$, go to Step 4d.

Step 6. [Assignment of a new row]

Set $C_i = r, P_r = P_r - \{l\}$; if $LR_r < 0$, set $U = U - \{r\}$, go to Step 4.

Otherwise, set $l = LR_r, P_r = P_r \cup \{l\}, r = LC_l$, repeat Step 6.

The efficiency of the algorithm is mainly due to the pointer technique utilized to locate the unexplored rows and the zero elements of the current cost matrix. It is worthwhile to note that, as far as storage requirement is concerned, for each set P_i ($i = 1, \dots, n$) only the pointer to the first element needs to be stored; in fact, if column j is contained in set P_i , the corresponding element of the current cost matrix (i.e., $a'_{i,j}$) is zero and it is thus possible to replace this element by the pointer to the column following j in set P_i .

We obtained a further improvement on the original Hungarian algorithm by modifying the choice of the initial solution, as described in Step 3 of the proposed algorithm.

Program

Fortran IV subroutine ASSCT, based on the algorithm previously presented, is completely self-contained and communication to it is made solely through the parameter list. The subroutine is called by means of the statement:

CALL ASSCT (N, A, C, T).

All the parameters are integer and their meanings are the following:

Input: N = number of rows and columns of the cost matrix;

A = cost matrix.

Output: C = assignment vector;

T = minimum assignment cost.

After execution of subroutine ASSCT, the values of the elements of the cost matrix are changed. Vector C must be dimensioned at least at N; matrix A at least at (N, N + 1). As presently dimensioned, the size limitation for ASSCT is N \leq 200.

Table I. Cost Range 1-100

<i>n</i>	A		B		C	
	Average	Maximum	Average	Maximum	Average	Maximum
50	0.16	0.20	0.21	0.31	0.23	0.33
100	0.54	0.66	0.93	1.15	1.10	1.27
150	0.96	1.10	2.10	2.81	2.79	3.42
200	1.40	1.68	4.37	5.27	5.38	7.24

Table II. Cost Range 1-1000

<i>n</i>	A		B		C	
	Average	Maximum	Average	Maximum	Average	Maximum
50	0.41	0.51	0.56	0.78	0.60	0.89
100	2.15	2.61	2.97	3.92	3.34	4.24
150	4.55	6.28	6.60	8.37	7.24	8.42
200	6.70	8.01	10.36	12.26	11.83	13.15

Table III. Cost Range 1-10,000

<i>n</i>	A		B		C	
	Average	Maximum	Average	Maximum	Average	Maximum
50	0.48	0.61	0.66	0.81	0.72	0.96
100	3.71	4.60	5.78	7.49	6.80	9.55
150	11.74	15.57	18.72	21.96	21.35	23.94
200	19.85	25.40	32.69	43.52	42.63	48.16

Table IV. Sparse Matrices, Cost Range 1-100, *n* = 200

Number of coefficients	1500	2250	3000	3750	4500
AP-AB	0.97	1.12	1.48	1.61	1.68
PD-AAL	1.63	1.14	1.89	1.29	1.80
SUPERT-2	1.26	1.57	1.98	2.17	2.53
ASSCT	15.29	5.80	6.49	2.62	2.17

Computational Results

Subroutine ASSCT was tested in a CDC 6600 with over 1500 random problems of varying sizes. No breakdown in the method occurred.

To evaluate the efficiency of the proposed algorithm, the computing times of subroutine ASSCT were compared with those of the most efficient algorithms for solution of the assignment problem for dense matrices [1, 3].

Tables I, II, and III show the computing times corresponding to the algorithms:

- A: Subroutine ASSCT.
- B: Hungarian algorithm as presented in [5] and coded in Fortran IV by the authors.
- C: Algorithm presented in [3] (the Fortran IV program, coded by Bourgeois and Lassalle, is taken from the CDC library of CERN (Geneva, Switzerland)).

In Tables I, II, and III the values of the cost matrix were generated as uniformly random integers in the ranges, respectively, 1-100, 1-1000, and 1-10,000. All codes were run on a CDC 6600. For each cost range, each algorithm, and each value of *n*, 20 different problems were solved, and the average and maximum computing times, expressed in seconds, are given.

The tables show that subroutine ASSCT is always superior to the other codes, mainly for large values of *n* and for small ranges of costs. In addition, for all codes the computing times get worse with the increase in size of coefficients; in fact, when the size increases, the number of zero elements of the current cost matrix decreases.

In order to evaluate the effect of sparseness on the performance of the proposed algorithm, the same test problems considered by Barr, Glover, and Klingman in [2] were solved on the same machine (a CDC-6600). Subroutine ASSCT was compared with the most efficient codes for the solution of sparse assignment problems, viz., codes AP-AB, PD-AAL, and SUPERT-2 presented, respectively, in [2], [6], and [1]; the corresponding computing times, expressed in seconds, are given in Table IV.¹

Table IV shows that the codes designed to solve sparse assignment problems are superior to subroutine ASSCT for very sparse matrices because this subroutine does not include any mechanism for taking advantage of sparsity. However, the trends of the computing times indicate that for fairly dense matrices the performance of the proposed algorithm greatly increases with respect to those of the other codes.

Further details of the algorithm and extensive computational results are given in [4].

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ALGORITHM

```

SUBROUTINE ASSCT ( N, A, C, T )                                10
  INTEGER A(200,201), C(200), CH(200), LC(200), LR(200),      20
  *   LZ(200), NZ(200), RH(201), SLC(200), SLR(200),          30
  *   U(201)                                                    40
  INTEGER H, Q, R, S, T                                        50
  EQUIVALENCE (LZ,RH), (NZ,CH)                                60
C                                                                 70
C THIS SUBROUTINE SOLVES THE SQUARE ASSIGNMENT PROBLEM        80
C THE MEANING OF THE INPUT PARAMETERS IS                       90
C N = NUMBER OF ROWS AND COLUMNS OF THE COST MATRIX, WITH   100
C   THE CURRENT DIMENSIONS THE MAXIMUM VALUE OF N IS 200     110
C A(I,J) = ELEMENT IN ROW I AND COLUMN J OF THE COST MATRIX  120
C ( AT THE END OF COMPUTATION THE ELEMENTS OF A ARE CHANGED) 130
C THE MEANING OF THE OUTPUT PARAMETERS IS                     140
C C(J) = ROW ASSIGNED TO COLUMN J (J=1,N)                     150
C T = COST OF THE OPTIMAL ASSIGNMENT                           160
C ALL PARAMETERS ARE INTEGER                                  170
C THE MEANING OF THE LOCAL VARIABLES IS                        180
C A(I,J) = ELEMENT OF THE COST MATRIX IF A(I,J) IS POSITIVE, 190
C   COLUMN OF THE UNASSIGNED ZERO FOLLOWING IN ROW I          1J
C   (I=1,N) THE UNASSIGNED ZERO OF COLUMN J (J=1,N)           210
C   IF A(I,J) IS NOT POSITIVE                                  220
C A(I,N+1) = COLUMN OF THE FIRST UNASSIGNED ZERO OF ROW I    230
C   (I=1,N)                                                    240
C CH(I) = COLUMN OF THE NEXT UNEXPLORED AND UNASSIGNED ZERO  250
C   OF ROW I (I=1,N)                                           260
C LC(J) = LABEL OF COLUMN J (J=1,N)                            270
C LR(I) = LABEL OF ROW I (I=1,N)                               280
C LZ(I) = COLUMN OF THE LAST UNASSIGNED ZERO OF ROW I (I=1,N) 290
C NZ(I) = COLUMN OF THE NEXT UNASSIGNED ZERO OF ROW I (I=1,N) 300
C RH(I) = UNEXPLORED ROW FOLLOWING THE UNEXPLORED ROW I      310
C   (I=1,N)                                                     320
C RH(N+1) = FIRST UNEXPLORED ROW                               330
C SLC(K) = K-TH ELEMENT CONTAINED IN THE SET OF THE LABELLED 340
C   COLUMNS                                                    350

```

¹ The times in Table IV were communicated to the authors by an anonymous referee who has given us permission to publish them.

130	C(J) = I	1130
	A(I,LJ) = A(I,J)	1140
	NZ(I) = -A(I,J)	1150
	LZ(I) = LJ	1160
	A(I,J) = 0	1170
140	CONTINUE	1180
C	RESEARCH OF A NEW ASSIGNMENT	1190
150	IF (U(NP1) .EQ. 0) RETURN	1200
	DO 160 I=1,N	1210
	CH(I) = 0	1220
	LC(I) = 0	1230
	LR(I) = 0	1240
	RH(I) = 0	1250
160	CONTINUE	1260
	RH(NP1) = -1	1270
	KSLC = 0	1280
	KSLR = 1	1290
	R = U(NP1)	1300
	LR(R) = -1	1310
	SLR(1) = R	1320
	IF (A(R,NP1) .EQ. 0) GO TO 220	1330
170	L = -A(R,NP1)	1340
	IF (A(R,L) .EQ. 0) GO TO 180	1350
	IF (RH(R) .NE. 0) GO TO 180	1360
	RH(R) = RH(NP1)	1370
	CH(R) = -A(R,L)	1380
	RH(NP1) = R	1390
180	IF (LC(L) .EQ. 0) GO TO 200	1400
	IF (RH(R) .EQ. 0) GO TO 210	1410
190	L = CH(R)	1420
	CH(R) = -A(R,L)	1430
	IF (A(R,L) .NE. 0) GO TO 180	1440
	RH(NP1) = RH(R)	1450
	RH(R) = 0	1460
	GO TO 180	1470
200	LC(L) = R	1480
	IF (C(L) .EQ. 0) GO TO 360	1490
	KSLC = KSLC+1	1500
	SLC(KSLC) = L	1510
	R = C(L)	1520
	LR(R) = L	1530
	KSLR = KSLR+1	1540
	SLR(KSLR) = R	1550
	IF (A(R,NP1) .NE. 0) GO TO 170	1560
210	CONTINUE	1570
	IF (RH(NP1) .GT. 0) GO TO 350	1580
C	REDUCTION OF THE CURRENT COST MATRIX	1590
220	H = MAXNUM	1600
	DO 240 J=1,N	1610
	IF (LC(J) .NE. 0) GO TO 240	1620
	DO 230 K=1,KSLR	1630
	I = SLR(K)	1640
	IF (A(I,J) .LT. H) H = A(I,J)	1650
230	CONTINUE	1660
240	CONTINUE	1670
	T = T+H	1680
	DO 290 J=1,N	1690
	IF (LC(J) .NE. 0) GO TO 290	1700
	DO 280 K=1,KSLR	1710
	I = SLR(K)	1720
	A(I,J) = A(I,J)-H	1730
	IF (A(I,J) .NE. 0) GO TO 280	1740
	IF (RH(I) .NE. 0) GO TO 250	1750
	RH(I) = RH(NP1)	1760
	CH(I) = J	1770
	RH(NP1) = I	1780
250	L = NP1	1790
260	NL = -A(I,L)	1800
	IF (NL .EQ. 0) GO TO 270	1810
	L = NL	1820
	GO TO 260	1830
270	A(I,L) = -J	1840
280	CONTINUE	1850
290	CONTINUE	1860
	IF (KSLC .EQ. 0) GO TO 350	1870
	DO 340 I=1,N	1880
	IF (LR(I) .NE. 0) GO TO 340	1890

DO 330 K=1,KSLC	1900
J = SLC(K)	1910
IF (A(I,J) .GT. 0) GO TO 320	1920
L = NP1	1930
300 NL = - A(I,L)	1940
IF (NL .EQ. J) GO TO 310	1950
L = NL	1960
GO TO 300	1970
310 A(I,L) = A(I,J)	1980
A(I,J) = H	1990
GO TO 330	2000
320 A(I,J) = A(I,J)+H	2010
330 CONTINUE	2020
340 CONTINUE	2030
350 R = RH(NP1)	2040
GO TO 190	2050
C ASSIGNMENT OF A NEW ROW	2060
360 C(L) = R	2070
M = NP1	2080
370 NM = -A(R,M)	2090
IF (NM .EQ. L) GO TO 380	2100
M = NM	2110
GO TO 370	2120
380 A(R,M) = A(R,L)	2130
A(R,L) = 0	2140
IF (LR(R) .LT. 0) GO TO 390	2150
L = LR(R)	2160
A(R,L) = A(R,NP1)	2170
A(R,NP1) = -L	2180
R = LC(L)	2190
GO TO 360	2200
390 U(NP1) = U(R)	2210
U(R) = 0	2220
GO TO 150	2230
END	2240

ALGORITHM 549

Weierstrass' Elliptic Functions [S21]

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Key Words and Phrases: Weierstrass' elliptic functions
CR Categories: 5.12
Language: Fortran

DESCRIPTION

Weierstrass' \mathcal{P} -function and its derivative can be approximated in the complex plane by a rational expression as described in [3, 4]. The following programs evaluate these approximations for the case of a unit period parallelogram. They are written according to the rules given in [2].

Programs PEQ and PEQ1 evaluate Weierstrass' \mathcal{P} -function in the equianharmonic case for a unit period parallelogram defined by the complex numbers

$$2\omega = \frac{1}{2} - \frac{1}{2}\sqrt{3}i, \quad 2\omega' = \frac{1}{2} + \frac{1}{2}\sqrt{3}i$$

(see [1, Sec. 18.13] for notation). Because of the periodicity of Weierstrass' functions it is sufficient to calculate them only in the fundamental rectangle of all $z = x + iy$ with

$$-\frac{1}{2} \leq x \leq \frac{1}{2}, \quad -\frac{\sqrt{3}}{4} \leq y \leq \frac{\sqrt{3}}{4}$$

which implies $|z| \leq \sqrt{7}/4$. The theoretical absolute error of the approximation in this rectangle can be estimated as in [3] by the following bounds:

$$\begin{array}{ll} \text{PEQ:} & 9.31 \times 10^{-19} \\ \text{PEQ1:} & 4.594 \times 10^{-18}. \end{array}$$

Numerical experiments were performed on different computers in order to obtain the actual errors of the programs. For large values of $|\mathcal{P}|$ the actual error is mainly due to the limited word length of the computer, whereas for small $|\mathcal{P}|$ the actual error approaches the theoretical error. Therefore a "mixed" error was calculated; it is the absolute error for $|\mathcal{P}| \leq 1$ and the relative error elsewhere. The bounds for the mixed errors were obtained by comparing the programs with ϑ -series evaluations for randomly and systematically chosen arguments. Each program was tested on the computers given in Table I for more than 20,000 argument values.

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Table I. Bounds for the Mixed Error for PEQ and PEQ1 on Different Computers

Computer	Mantissa length			
	Bits	Decimal	PEQ	PEQ1
IBM 370/168 (double precision)	56	16	2×10^{-15}	2×10^{-14}
CDC Cyber 172	48	14	7×10^{-13}	7×10^{-13}
TR 440	38	11	2×10^{-9}	3×10^{-9}

Table II. Bounds for the Mixed Error for PLEM and PLEM1

Computer	PLEM	PLEM1
IBM 370/168 (double precision)	4×10^{-15}	8×10^{-15}
CDC Cyber 172	6×10^{-13}	6×10^{-13}
TR 440	4×10^{-9}	5×10^{-9}
Theoretical absolute error	4.482×10^{-17}	8.033×10^{-17}

The difference between the numerical accuracy obtainable by using full word length of the computer and the values given in Table I is due to

- (1) a loss of nearly one significant figure in computing powers of the variable z ;
- (2) a loss of one significant figure in evaluating the first line of the last assignment statement in each program.

Programs PLEM and PLEM1 evaluate Weierstrass' \mathcal{P} -function and its derivative in the lemniscatic case where

$$2\omega = 1, \quad 2\omega' = i.$$

The fundamental rectangle is defined by

$$-\frac{1}{2} \leq x \leq \frac{1}{2}, \quad -\frac{1}{2} \leq y \leq \frac{1}{2}$$

so that $|z| \leq 1/\sqrt{2}$. The theoretical absolute error as given in [4] and the actual mixed error which was found by numerical experiments are given in Table II.

In using the programs it should be observed that the functions have poles at the lattice points $2\omega M + 2\omega' N$. If z happens to coincide with one of these poles, an overflow occurs. This can be avoided by checking after the reduction of z (see comments in the program listings) whether the reduced value of z is close to zero.

Full machine accuracy can be obtained in all programs by specifying the argument z and all its powers in greater precision and by performing all operations of the first line of the last assignment statement in double precision.

Note. The programs are written in single precision. On IBM computers they should be declared as

$$\text{COMPLEX FUNCTION } \left\{ \begin{array}{l} \text{PEQ} \\ \text{PEQ1} \\ \text{PLEM} \\ \text{PLEM1} \end{array} \right\} * 16(Z).$$

The first declaration should be `COMPLEX * 16` and the second `REAL * 8`. All constants and functions should be given in double precision.

ACKNOWLEDGMENTS

The programs were written and tested on the IBM 370/168 of the Central Institute for Applied Mathematics of the Nuclear Research Center Jülich. They were also tested on the CDC Cyber 172 and the Telefunken TR 440 of the

Regionales Rechenzentrum Erlangen. The author is very much indebted to W. Alt of the University of Bayreuth for kind assistance.

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ALGORITHM

```

COMPLEX Z, PLEM, PLEM1, CR, CI, P, P1          MAN 10
REAL W2, LC, ALPHA, FP, FP1, EPS, EPS1, A      MAN 20
C                                               MAN 30
C *****MAN 40
C                                               MAN 50
C THE PROGRAMS FOR WEIERSTRASS: P-FUNCTION AND ITS DERIVATIVE IN THE MAN 60
C LEMNISCATIC CASE ARE TESTED FOR THE ARGUMENT VALUES          MAN 70
C Z = 0.25, 0.5, 0.5 + 0.5*I, 0.25 + 0.25*I          MAN 80
C AND 10.5 - 7.5*I          MAN 90
C (SEE M. ABRAMOWITZ AND I. A. STEGUN & HANDBOOK OF MATHEMATICAL MAN 100
C FUNCTIONS, 18.14 FOR CORRESPONDING FUNCTION VALUES)          MAN 110
C THE DRIVER PROGRAM PRINTS THE MAXIMAL RELATIVE ERRORS OF P AND P: MAN 120
C                                               MAN 130
C *****MAN 140
C                                               MAN 150
C                                               MAN 160
C EVALUATION OF CONSTANTS          MAN 170
C =====          MAN 180
C                                               MAN 190
C LC = 2.6220575542921198104648396          MAN 200
C                                               MAN 210
C LC IS THE LEMNISCATE CONSTANT          MAN 220
C                                               MAN 230
C W2 = 2.0*LC/SQRT(2.0)          MAN 240
C ALPHA = 1.0 + SQRT(2.0)          MAN 250
C CR = (1.0,0.0)          MAN 260
C CI = (0.0,1.0)          MAN 270
C FP = 1.0/(W2*W2)          MAN 280
C FP1 = FP/W2          MAN 290
C                                               MAN 300
C * * * * *          MAN 310
C                                               MAN 320
C Z = (0.25,0.0)          MAN 330
C P = PLEM(Z)          MAN 340
C P1 = PLEM1(Z)          MAN 350
C EPS = CABS(P*FP-0.5*ALPHA*CR)/CABS(P)          MAN 360
C EPS1 = CABS(P1*FP1+ALPHA*CR)/CABS(P1)          MAN 370
C                                               MAN 380
C * * * * *          MAN 390
C                                               MAN 400
C Z = (0.5,0.0)          MAN 410
C P = PLEM(Z)          MAN 420
C P1 = PLEM1(Z)          MAN 430
C A = CABS(P*FP-0.5*CR)/CABS(P)          MAN 440
C IF (A.GT.EPS) EPS = A          MAN 450
C A = CABS(P1*FP1)          MAN 460
C IF (A.GT.EPS1) EPS1 = A          MAN 470
C                                               MAN 480
C * * * * *          MAN 490
C                                               MAN 500
C Z = (0.5,0.5)          MAN 510
C P = PLEM(Z)          MAN 520
C P1 = PLEM1(Z)          MAN 530
C A = CABS(P*FP)          MAN 540
C IF (A.GT.EPS) EPS = A          MAN 550
C A = CABS(P1*FP1)          MAN 560
C IF (A.GT.EPS1) EPS1 = A          MAN 570

```

C		MAN	580
C	* * * * *	MAN	590
C		MAN	600
	Z = (0.25,0.25)	MAN	610
	P = PLEM(Z)	MAN	620
	P1 = PLEM1(Z)	MAN	630
	A = CABS(P*FP+0.5*CI)/CABS(P)	MAN	640
	IF (A.GT.EPS) EPS = A	MAN	650
	A = CABS(P1*FP1-(CR+CI)/SQRT(2.0))/CABS(P1)	MAN	660
	IF (A.GT.EPS1) EPS1 = A	MAN	670
C		MAN	680
C	* * * * *	MAN	690
C		MAN	700
	Z = (10.5,-7.5)	MAN	710
C		MAN	720
C	THIS Z IS FOR TESTING THE REDUCTION TO FUNDAMENTAL PARALLELOGRAM.	MAN	730
C	BECAUSE OF THIS REDUCTION, THE ERROR MIGHT BE AN ORDER OF	MAN	740
C	MAGNITUDE GREATER THAN THE BOUND GIVEN IN THE TEXT	MAN	750
C		MAN	760
	P = PLEM(Z)	MAN	770
	P1 = PLEM1(Z)	MAN	780
	A = CABS(P*FP)	MAN	790
	IF (A.GT.EPS) EPS = A	MAN	800
	A = CABS(P1*FP1)	MAN	810
	IF (A.GT.EPS1) EPS1 = A	MAN	820
C		MAN	830
C	* * * * *	MAN	840
C	* * * * *	MAN	850
C		MAN	860
	WRITE (6,99999)	MAN	870
	WRITE (6,99998) EPS	MAN	880
	WRITE (6,99997) EPS1	MAN	890
	STOP	MAN	900
99999	FORMAT (51H1 MAXIMAL RELATIVE ERROR FOR WEIERSTRASS P-FUNCTI,	MAN	910
	* 2HON/49H IN THE LEMNISCATIC CASE AT FIVE SAMPLE POINTS/4X,	MAN	920
	* 50(1H=))	MAN	930
99998	FORMAT (///25H RELATIVE ERROR FOR P=, 5X, E10.3)	MAN	940
99997	FORMAT (//26H RELATIVE ERROR FOR P1=, 4X, E10.3)	MAN	950
	END	MAN	960
	COMPLEX FUNCTION PEQ(Z)	PEQ	10
C		PEQ	20
C	WEIERSTRASS: P-FUNCTION IN THE EQUIANHARMONIC CASE	PEQ	30
C	FOR COMPLEX ARGUMENT WITH UNIT PERIOD PARALLELOGRAM	PEQ	40
C		PEQ	50
	COMPLEX Z, Z2, Z4, Z6	PEQ	60
	REAL ZR, ZI	PEQ	70
	INTEGER M, N	PEQ	80
C		PEQ	90
C	REDUCTION TO FUNDAMENTAL PARALLELOGRAM	PEQ	100
C		PEQ	110
	ZI = 1.1547005383792515E0*AIMAG(Z) + 0.5E0	PEQ	120
	M = INT(ZI)	PEQ	130
	IF (ZI.LT.0E0) M = M - 1	PEQ	140
	ZR = REAL(Z) - 0.5E0*FLOAT(M) + 0.5E0	PEQ	150
	N = INT(ZR)	PEQ	160
	IF (ZR.LT.0E0) N = N - 1	PEQ	170
	Z2 = Z - FLOAT(N) - (0.5E0,0.86602540378443865E0)*FLOAT(M)	PEQ	180
C		PEQ	190
C	IF Z2=0 THEN Z COINCIDES WITH A LATTICE POINT.	PEQ	200
C	SINCE P HAS POLES AT THE LATTICE POINTS,	PEQ	210
C	A DIVISION ERROR WILL OCCUR	PEQ	220
C		PEQ	230
	Z2 = Z2*Z2	PEQ	240
	Z4 = Z2*Z2	PEQ	250
	Z6 = Z4*Z2	PEQ	260
	PEQ = 1E0/Z2 + 6E0*Z4*(5E0+Z6)/(1E0-Z6)**2 + Z4*	PEQ	270
	* ((((-2.6427662E-10*Z6+1.610954818E-8)*Z6+7.38610752879E-6)*	PEQ	280
	* Z6+4.3991444671178E-4)*Z6+7.477288220490697E-2)*	PEQ	290
	* Z6-6.8484153287299201E-1)/((((6.2252191E-10*Z6+2.553314573E-7)*	PEQ	300
	* Z6-2.619832920421E-5)*Z6-5.6444801847646E-4)*	PEQ	310
	* Z6+4.565553484820106E-2)*Z6+1E0)	PEQ	320
	RETURN	PEQ	330
	END	PEQ	340

	COMPLEX FUNCTION PLEMI(Z)	PLE	10
C		PLE	20
C	FIRST DERIVATIVE OF WEIERSTRASS: P-FUNCTION IN THE	PLE	30
C	LEMNISCATIC CASE FOR COMPLEX ARGUMENT	PLE	40
C	WITH UNIT PERIOD PARALLELOGRAM	PLE	50
C		PLE	60
	COMPLEX Z, Z1, Z3, Z4	PLE	70
	REAL ZR, ZI	PLE	80
	INTEGER M, N	PLE	90
C		PLE	100
C	REDUCTION TO FUNDAMENTAL PARALLELOGRAM	PLE	110
C		PLE	120
	ZR = REAL(Z) + 0.5E0	PLE	130
	ZI = ALMAG(Z) + 0.5E0	PLE	140
	M = INT(ZR)	PLE	150
	N = INT(ZI)	PLE	160
	IF (ZR.LT.0E0) M = M - 1	PLE	170
	IF (ZI.LT.0E0) N = N - 1	PLE	180
	Z1 = Z - FLOAT(M) - (0E0,1E0)*FLOAT(N)	PLE	190
C		PLE	200
C	IF Z1=0 THEN Z COINCIDES WITH A LATTICE POINT.	PLE	210
C	SINCE P: HAS POLES AT THE LATTICE POINTS,	PLE	220
C	A DIVISION ERROR WILL OCCUR	PLE	230
C		PLE	240
	Z3 = Z1*Z1*Z1	PLE	250
	Z4 = Z3*Z1	PLE	260
	PLEMI = (((1E1*Z4+9E1)*Z4+3E1)*Z4-2E0)/(Z1*(1E0-Z4))**3 +	PLE	270
	* Z1*(((((((-3.9046302E-9*Z4-1.001487137E-8)*Z4+5.9573043092E-7)	PLE	280
	* *Z4-2.482518130524E-5)*Z4+1.4557266595395E-4)*	PLE	290
	* Z4+4.56633655643206E-3)*Z4+6.224782572111135E-2)*	PLE	300
	* Z4+1.038527937794269E-2)*Z4+1.19804620802637942E0)*	PLE	310
	* Z4+6.42791439683811718E0)*Z4-5.09272798707661477E0)/	PLE	320
	* ((((((((((4.726888E-11*Z4-3.0667983E-9)*Z4+1.0087596089E-7)*	PLE	330
	* Z4-8.0606833451E-8)*Z4+1.184299251664E-5)*Z4-2.3096723361547E-4)*	PLE	340
	* Z4-2.90730903142055E-3)*Z4+1.338392411135511E-2)*	PLE	350
	* Z4+2.3098639320021426E-1)*Z4+8.4719880964554148E-1)*Z4+1E0)	PLE	360
	RETURN	PLE	370
	END	PLE	380
	COMPLEX Z, PEQ, PEQ1, CR, CI, P, P1	MAN	10
	REAL W2, SQ3, FP, FP1, EPS, EPS1, A, C	MAN	20
C		MAN	30
C	*****MAN	MAN	40
C		MAN	50
C	THE PROGRAMS FOR WEIERSTRASS: P-FUNCTION AND ITS DERIVATIVE IN THE	MAN	60
C	EQUIANHARMONIC CASE ARE TESTED FOR THE ARGUMENT VALUES	MAN	70
C	Z = 0.25, 0.5, 0.5 + 0.5/SQRT(3)*I, 0.25 + 0.25/SQRT(3)*I	MAN	80
C	AND 2.5 - 47/6*SQRT(3)*I	MAN	90
C	(SEE M. ABRAMOWITZ AND I. A. STEGUN& HANDBOOK OF MATHEMATICAL	MAN	100
C	FUNCTIONS, 18.13. FOR CORRESPONDING FUNCTION VALUES)	MAN	110
C	THE DRIVER PROGRAM PRINTS THE MAXIMAL RELATIVE ERRORS OF P AND P:	MAN	120
C		MAN	130
C	*****MAN	MAN	140
C		MAN	150
C		MAN	160
C	EVALUATION OF CONSTANTS	MAN	170
C		MAN	180
C	=====	MAN	190
C		MAN	200
	SQ3 = SQRT(3.0)	MAN	210
	CR = (1.0,0.0)	MAN	220
	CI = (0.0,1.0)	MAN	230
	C = 4.0**(-1.0/3.0)	MAN	240
	W2 = 1.52995403705719287491319417231	MAN	250
	W2 = 2.0*W2	MAN	260
	FP = 1.0/(W2*W2)	MAN	270
	FP1 = FP/W2	MAN	280
C		MAN	290
C	* * * * *	MAN	300
C		MAN	310
	Z = (0.25,0.0)	MAN	320
	P = PEQ(Z)	MAN	330
	P1 = PEQ1(Z)	MAN	340
	EPS = CABS(P*FP-C*(1.0+SQ3)*CR)/CABS(P)	MAN	350
	EPS1 = CABS(P1*FP1+3.0**0.75*SQRT(2.0+SQ3)*CR)/CABS(P1)	MAN	360

C		MAN	370
C	* * * * *	MAN	380
C		MAN	390
	Z = (0.5, 0.0)	MAN	400
	P = PEQ(Z)	MAN	410
	P1 = PEQ1(Z)	MAN	420
	A = CABS(P*FP-C*CR)/CABS(P)	MAN	430
	IF (A.GT.EPS) EPS = A	MAN	440
	A = CABS(P1*FP1)	MAN	450
	IF (A.GT.EPS1) EPS1 = A	MAN	460
C		MAN	470
C	* * * * *	MAN	480
C		MAN	490
	Z = 0.5*CR + 0.5/SQ3*CI	MAN	500
	P = PEQ(Z)	MAN	510
	P1 = PEQ1(Z)	MAN	520
	A = CABS(P*FP)	MAN	530
	IF (A.GT.EPS) EPS = A	MAN	540
	A = CABS(P1*FP1-CI)/CABS(P1)	MAN	550
	IF (A.GT.EPS1) EPS1 = A	MAN	560
C		MAN	570
C	* * * * *	MAN	580
C		MAN	590
	Z = Z*0.5	MAN	600
	P = PEQ(Z)	MAN	610
	P1 = PEQ1(Z)	MAN	620
	A = CABS(P*FP+2.0**(1.0/3.0)*0.5*(-CR+SQ3*CI))/CABS(P)	MAN	630
	IF (A.GT.EPS) EPS = A	MAN	640
	A = CABS(P1*FP1-(0.0, 3.0))/CABS(P1)	MAN	650
	IF (A.GT.EPS1) EPS1 = A	MAN	660
C		MAN	670
C	* * * * *	MAN	680
C		MAN	690
	Z = 2.5*CR - 47.0*SQ3/6.0*CI	MAN	700
C		MAN	710
C	THIS Z IS FOR TESTING THE REDUCTION TO FUNDAMENTAL PARALLELOGRAM.	MAN	720
C	BECAUSE OF THIS REDUCTION, THE ERROR MIGHT BE AN ORDER OF	MAN	730
C	MAGNITUDE GREATER THAN THE BOUND GIVEN IN THE TEXT	MAN	740
C		MAN	750
	P = PEQ(Z)	MAN	760
	P1 = PEQ1(Z)	MAN	770
	A = CABS(P*FP)	MAN	780
	IF (A.GT.EPS) EPS = A	MAN	790
	A = CABS(P1*FP1-CI)/CABS(P1)	MAN	800
	IF (A.GT.EPS1) EPS1 = A	MAN	810
C		MAN	820
C	* * * * *	MAN	830
C	* * * * *	MAN	840
C		MAN	850
	WRITE (6,99999)	MAN	860
	WRITE (6,99998) EPS	MAN	870
	WRITE (6,99997) EPS1	MAN	880
	STOP	MAN	890
99999	FORMAT (51H1 MAXIMAL RELATIVE ERROR FOR WEIERSTRASS P-FUNCTI,	MAN	900
	* 2HON/52H IN THE EQUIANHARMONIC CASE AT FIVE SAMPLE POINTS/4X,	MAN	910
	* 50(1H=)	MAN	920
99998	FORMAT (///25H RELATIVE ERROR FOR P=, 5X, E10.3)	MAN	930
99997	FORMAT (//26H RELATIVE ERROR FOR P1=, 4X, E10.3)	MAN	940
	END	MAN	950

ALGORITHM 550

Solid Polyhedron Measures [Z]

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General Dynamics

and

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Key Words and Phrases: polyhedron, graphics, numerical integration

CR Categories: 3.2, 5.16, 8.2

Language: Fortran

DESCRIPTION

Surface area, centroid, volume, weight, moments, and products of inertia are computed by numerical integration over the bounding surfaces of solid polyhedra. This routine operates directly on geometric definitions consisting of vertex coordinates and face lists, in a format similar to the format employed in graphics and in plotting computer programs [1, 3, 5]. Integration is performed through a quadrature rule that is exact for the equations involved.

METHOD

The several equations, eqs. (1) through (10), to be integrated over the surface of the polyhedra were derived from volume integral forms through application of the Gauss divergence theorem.

$$\text{mass} = \frac{\rho}{3} \iint [xn_x + yn_y + zn_z] dS = M \quad (1)$$

$$CG_x = \frac{\rho}{2M} \iint [x^2n_x] dS \quad (2)$$

$$CG_y = \frac{\rho}{2M} \iint [y^2n_y] dS \quad (3)$$

$$CG_z = \frac{\rho}{2M} \iint [z^2n_z] dS \quad (4)$$

$$I_{xx} = \frac{\rho}{3} \iint [y^3n_y + z^3n_z] dS \quad (5)$$

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$$I_{yy} = \frac{\rho}{3} \iint [z^3 n_z + x^3 n_x] dS \quad (6)$$

$$I_{zz} = \frac{\rho}{3} \iint [x^3 n_x + y^3 n_y] dS \quad (7)$$

$$P_{xy} = \frac{-\rho}{2} \iint [yx^2 n_x] dS \quad (8)$$

$$P_{yz} = \frac{-\rho}{2} \iint [zy^2 n_y] dS \quad (9)$$

$$P_{zx} = \frac{-\rho}{2} \iint [xz^2 n_z] dS \quad (10)$$

where $\iint dS$ implies integration over the entire bounding surface; x , y , and z are right-hand Cartesian coordinates; n_x , n_y , and n_z are components of the outward directed normal vector of unit length; ρ is density per cubic unit; CG_x , CG_y , and CG_z are coordinates of the centroid; I_{xx} , I_{yy} , and I_{zz} are moments of inertia about the x , y , and z axes; P_{xy} , P_{yz} , and P_{zx} are products of inertia with respect to the subscripted axes with sign convention as shown.

Consider, for example, the moment of inertia about the x axis (as given in [2]):

$$I_{xx} = \iiint \rho(y^2 + z^2) dV.$$

We define a specific vector function \mathbf{F} as follows:

$$\mathbf{F} = \frac{\rho}{3} \{y^3 \mathbf{j} + z^3 \mathbf{k}\}$$

where \mathbf{j} and \mathbf{k} are unit vectors in the coordinate directions y and z , respectively. The divergence of \mathbf{F} is

$$\text{div } \mathbf{F} = \frac{\rho}{3} \left\{ \frac{\partial(y^3)}{\partial y} + \frac{\partial(z^3)}{\partial z} \right\} = \rho(y^2 + z^2).$$

Gauss's divergence theorem states

$$\iiint \text{div } \mathbf{F} dV = \iint \mathbf{F} \cdot \mathbf{n} dS$$

yielding

$$\iiint \rho(y^2 + z^2) dV = \iint \frac{\rho}{3} (y^3 \mathbf{ny} + z^3 \mathbf{nz}) dS$$

which verifies that the surface integral form, eq. (5), is the equivalent of the volume representation.

Computation of the surface integrals is done through a quadrature rule applied to the triangular elements of the surface polygons. We employ a four-point rule, eq. (11), that is exact for polynomials up to degree three. This rule, a minimum-point exact rule for our purpose [6] converts the continuous integral into a sum of four discrete values evaluated at specific points in each triangle. Point locations and associated weights are as follows:

$$\iint f dA = \text{area}[f(p_1)w_1 + f(p_2)w_2 + f(p_3)w_3 + f(p_4)w_4] \quad (11)$$

where p_1 is the centroid of the triangle, $w_1 = -27/48$, $w_2 = w_3 = w_4 = 25/48$. p_2 , p_3 , and p_4 are located along lines from the centroid to the triangle vertices such that the distance to these points is equal to 40 percent of the distance from the centroid to the respective vertex.

Each surface polygon is treated as a group of triangles that are constructed as follows: Vertices 1, 2, and 3 form the first triangle, vertices 1, 3, and 4, form the second triangle, etc., with 1, $n - 1$, and n being the last triangle in a perimeter of n vertices. Normal vectors are computed for each triangle to ensure that concave polygons yield negative contributions when the perimeter reverses direction relative to the initial vertex. Contributions to mass, CG, products of inertia, and moments of inertia are accumulated in sequence to cascade the calculations and thereby reduce arithmetic operations. Global factors such as density are applied after all faces have been processed. Properties are finally shifted to parallel axes through the computed centroid. In this form, standard methods can be employed to yield values about arbitrary axes or to solve for principal values.

DESCRIPTION OF THE ALGORITHM

The algorithm is presented as two subroutines: PROPS and SRFINT, which are structured to operate on polyhedrons defined as a series of vertices and associated connectivity (faces) lists. Specific details of an appropriate data structure are illustrated by the example shown in Figure 1 and a driver program based on this format is included. Vertices are numbered, and defined in terms of their components in a right-hand rectangular Cartesian coordinate system. Faces are also numbered, and listed as a series of vertices in the order encountered while traversing the perimeters. Perimeters are traversed in a direction that is consistent with an inward surface vector definition. This ordering convention can be envisioned as follows: If one were to stand on the face with the left foot on the edge of the face and the right foot on the interior of the face and then traverse forward around the perimeter, vertices would be encountered in the required order. Perimeter lists are closed (the first vertex number is repeated at the end of the list) and each face consists of one or more perimeters. These aspects are illustrated by the example shown in Figure 1. Large faces are accommodated for in the driver program by continuing perimeter lists on subsequent cards using all data fields (including the first field normally used for the face number). A repeat of the initial vertex designates the end of the list.

Errors in the input data can occur, and are often difficult to detect unless systematic tests of data integrity are adopted. Geometric tests can be made to insure that the polyhedron is connected and has planar faces. These tests will vary with the application; nonplanar faces, for example, may be permitted in a mechanically deformed solid, relying on the triangulation process within the algorithm to enforce planarity. In our experience the most useful test has been a connectivity evaluation that sorts all edges and proceeds to match each edge against its reciprocal (same edge lists in the opposite order on another face). This test is easily accomplished while generating plots [3]. If each edge has one and only one reciprocal in the face lists there is reasonable assurance that the connectivity is valid. Another data error encountered in extensive polygon definitions is the omission of coordinates of one or more vertices. This error will usually result in coordinate values of zero, which can be tested, and appropriate cautionary messages printed out to aid the correction process. Large data sets may require graphic-generation techniques such as the use of digitizers or interactive entry devices. In most instances computer plots should be obtained as a further data check. Our experience is primarily with computations relating to designed components which are defined by engineering drawings. These drawings are reduced to appropriate card formats through digitizer equipment interfaced to a keypunch machine. Perspective plots (Figures 2 and 3) generated through a derivative of the Loutrel Algorithm [3] have been valuable guides to data integrity.

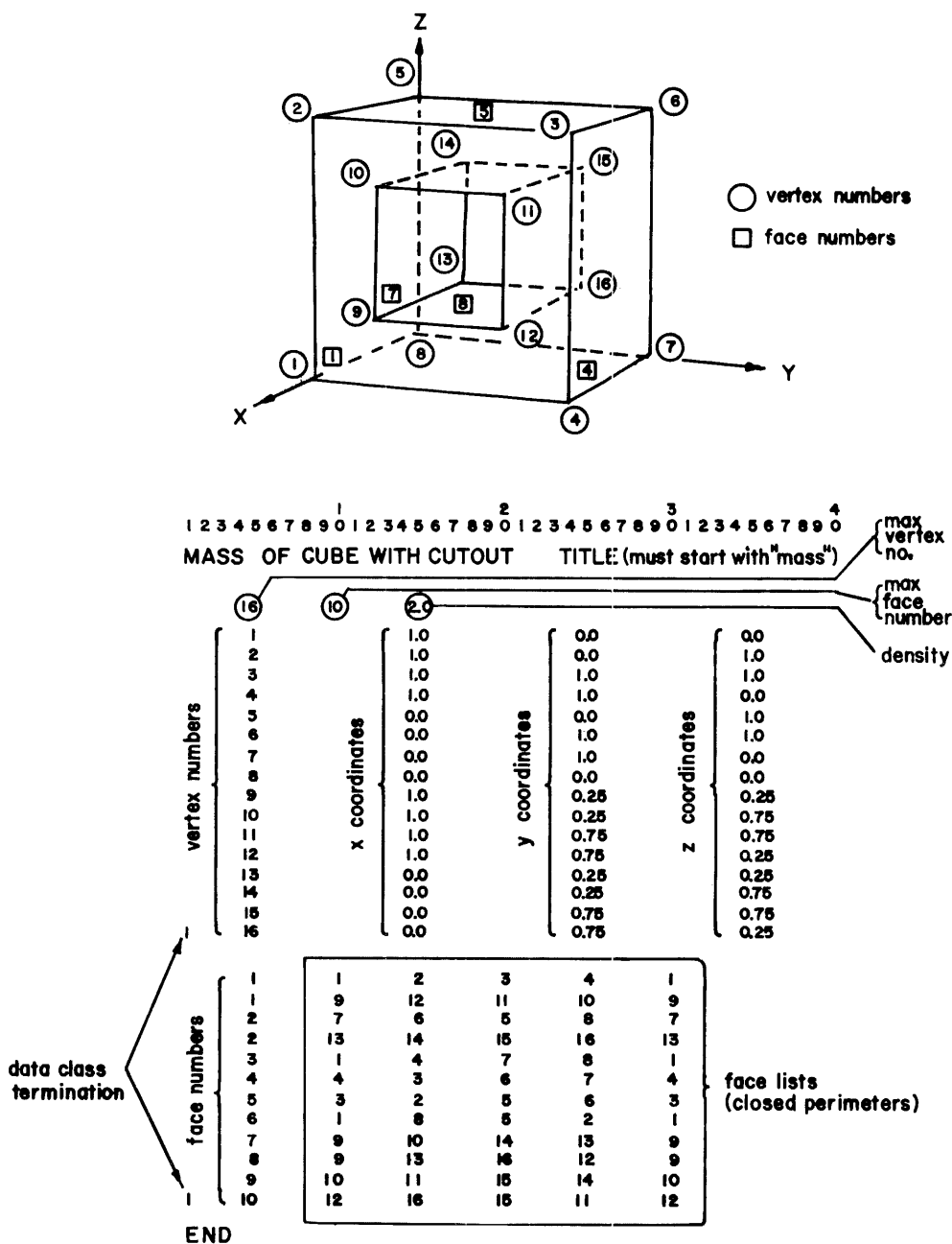


Fig. 1. Input definitions.

TESTS

This algorithm has been in use for several years and checked for simple geometries such as cubes and prisms as well as for complex configurations. Comparisons have also been made with fabricated components where computed values were compared with measured quantities of finished parts. Accuracies in all instances were within manufacturing tolerances. Computational errors arise from roundoff operations within the computer and can be minimized by multiple-precision arithmetic or through specific protection algorithms [4].

Several specific test problems are presented to illustrate the accuracies obtainable in single and double precision with the IBM 370 equipment. The first

Table I. Solutions Using Single and Double Precision with IBM 370 Computer

<i>Single-Precision Solution to Test Problem 1 (100.0-inch cube with cutout 0.01-inch wall.)</i>			
Density	2.00000		
Area	79999.8750 (79999.9992) ^a		
Volume	399.7500 (399.96) ^a		
Weight	799.50000 (799.92) ^a		
	Axes		
	<i>x</i>	<i>y</i>	<i>z</i>
CG			
coordinates	50.031	50.071	49.871
Moment	(50.0) ^a	(50.0) ^a	(50.0) ^a
from origin	6589440.00	5943488.00	5953024.00
from CG	2596525.00 (2665866.666) ^a	1953777.00 (1999533.36) ^a	1947323.00 (1999533.36) ^a
	Axes		
	<i>xy</i>	<i>yz</i>	<i>zx</i>
Product moment			
from origin	1992960.00	1988096.00	1986816.00
from CG	-9891.000 (0.0) ^a	-8345.000 (0.0) ^a	-8029.000 (0.0) ^a
<i>Double-Precision Solution to Test Problem 1 (100.0-inch cube with cutout 0.01-inch wall.)</i>			
Density	2.00000		
Area	79999.99920		
Volume	399.95994		
Weight	799.91988		
	Axes		
	<i>x</i>	<i>y</i>	<i>z</i>
CG			
coordinates	50.00000	50.00000	50.00000
Moment			
from origin	6665464.73061	5999131.51563	5999131.51563
from CG	2665866.08647	1999532.87149	1999532.87149
	Axes		
	<i>xy</i>	<i>yz</i>	<i>zx</i>
Product moment			
from origin	1999799.32206	1999799.32206	1999799.32206
from CG	-0.000000	-0.00001	-0.00001

^a True values shown below analogous computer results.

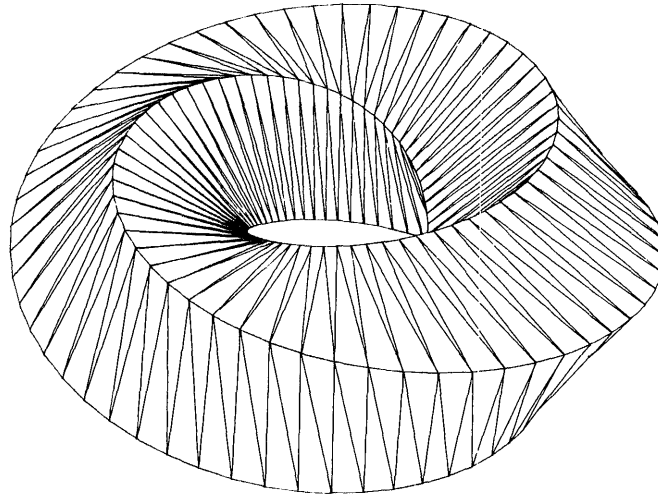


Fig. 2. Test Problem 2 generated solid.

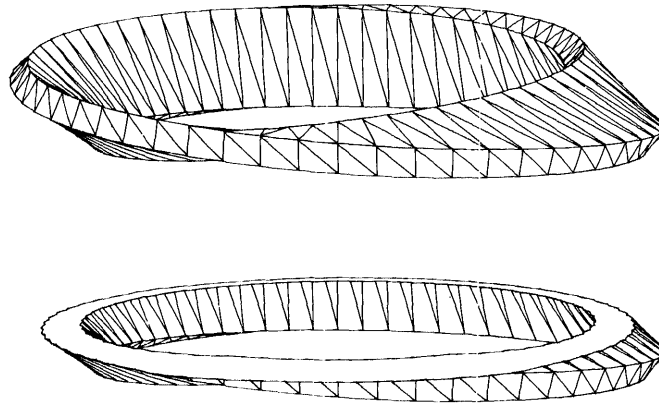


Fig. 3. Test Problem 3 Möbius and sectioned Möbius.

problem consists of a 100.0-inch cube, with cutout as in Figure 1. Cutout dimensions are such that the remaining wall thickness is 0.01 inch. Computations involve powers of coordinates which differ by the wall thickness; therefore, as the wall gets sufficiently thin, accuracy deteriorates. The complete single- and double-precision solution to this problem is presented in Table I. It can be seen that inaccuracies are evident in the single-precision solution, although even this problem, which is somewhat extreme for practical purposes, did not produce catastrophic errors. A second test has been constructed to indicate the accuracies related to geometries containing numerous faces with diverse orientations. This second test relates to the symmetry of results obtained for a generated configuration depicted in Figure 2. It was constructed by displacing an 8-inch square cross section normal to its plane in such a manner that its center traversed a 10-inch-radius circle while rotating 180 degrees in its plane during one traverse of the circle. The circle has its center at the origin and lies in the x - y plane. Single- and double-precision solutions are presented in Table II. Moments about x and y should be equal and the centroid and products of inertia components should be zero. It is of interest to note that the volume is some 2 percent smaller than would be obtained with a continuous solid generated in this manner. This discrepancy arises from the inscribing character of the polyhedron construction and would decrease as the number of faces increase. A final test is presented in Figure 3.

Table II. Computer Solutions to Example Problem 2

<i>Single-Precision Solution to Problem 2</i>			
Density	2.00000		
Area	2058.57715		
Volume	3941.82104		
Weight	7883.64062		
		Axes	
		<i>x</i>	<i>y</i>
		<i>z</i>	
CG			
coordinates	0.000	-0.000	0.000
Moment			
from origin	497116.062	497103.187	911423.062
from CG	497116.062	497103.187	911423.062
		Axes	
		<i>xy</i>	<i>yz</i>
		<i>zx</i>	
Product moment			
from origin	-23.484	-0.063	-0.033
from CG	-23.484	-0.063	-0.033
<i>Double-Precision Solution to Problem 2</i>			
Density	2.00000		
Area	2058.63206		
Volume	3942.03508		
Weight	7884.07016		
		Axes	
		<i>x</i>	<i>y</i>
		<i>z</i>	
CG			
coordinates	-0.00000	0.00000	0.00000
Moment			
from origin	497209.13864	497209.12703	911523.19783
from CG	497209.13864	497209.12703	911523.19783
		Axes	
		<i>xy</i>	<i>yz</i>
		<i>zx</i>	
Product moment			
from origin	-0.00155	0.00022	-0.00010
from CG	-0.00155	0.00022	-0.00010

Table III

<i>Solution to Complete Solid Test Problem 3</i>			
Density	2.00000		
Area	630.20361		
Volume	241.44048		
Weight	482.88086		
	Axes		
	<i>x</i>	<i>y</i>	<i>z</i>
CG			
coordinates	0.03247	-0.00031	0.00000
Moment			
from origin	24961.5117	24960.8477	49240.2383
from CG	24961.5117	24960.3359	49239.7266
	Axes		
	<i>xy</i>	<i>yz</i>	<i>zx</i>
Product moment			
from origin	-0.21539	-309.58984	0.00619
from CG	-0.21049	-309.58984	0.00618
<i>Solution to Sectioned Solid Test Problem 3</i>			
Density	2.00000		
Area	424.16016		
Volume	120.71504		
Weight	241.43008		
	Axes		
	<i>x</i>	<i>y</i>	<i>z</i>
CG			
coordinates	0.03247	-0.75071	0.67662
Moment			
from origin	12478.8516	12479.5312	24617.2148
from CG	12232.2578	12368.7422	24480.8984
	Axes		
	<i>xy</i>	<i>yz</i>	<i>zx</i>
Product moment			
from origin	-375.08545	-154.91779	-445.60889
from CG	-369.20068	-32.28427	-450.91260

This third test problem is similar in configuration to the second except that the cross section is now a 1- by 4-inch rectangle (finite thickness Möbius strip). A plane through the generating circle of this configuration bisects it identically. Solutions in single precision are presented to the complete and the sectioned solid in Table III where it can be seen that the volume, weight, moments about the axes, and the y - z product moment are halved as required. The x coordinate of the centroid is not zero due to the discretization errors associated with the geometry construction.

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ALGORITHM

```

C PURPOSE
C TO CALCULATE MASS PROPERTIES OF GENERAL THREE
C DIMENSIONAL SOLIDS.
C
C DESCRIPTION
C THE SOLID IS DESCRIBED BY A SET OF SURFACES WHICH ARE
C TRIANGULATED, AND A FOUR POINT GAUSSIAN QUADRATURE OVER
C THE TRIANGLES IS USED IN CALCULATING SURFACE AREA,
C VOLUME, AND MASS PROPERTIES.
C
C
C COMMON MAXS,MAXV,
. AREA,WEIGHT,VOL,DENS,
. CG(3),IO(3),ICG(3),PR(3),PRCG(3),
. ERROR,NERR,
. HED(80),IN,OT,KT,
. N8,MAXKV
LOGICAL ERROR
INTEGER OT
INTEGER TITLE,END,HED
DIMENSION A(10000),KA(6000)
REAL IO,ICG
DATA TITLE,END /1HM,1HE/
C
C THE FOLLOWING TWO STATEMENTS SET STORAGE LIMITS.
C DIMENSION A(10000),KA(6000)
MAXKV=6000
IN=5
OT=6
C
C TITLE AND CONTROL CARDS ARE READ FOR CURRENT CASE.
C
1 READ (IN,4) HED
IF (HED(1) .EQ. TITLE) GO TO 2
IF (HED(1) .NE. END ) GO TO 1
C END OF JOB
STOP
C
C ARRAY POINTERS ARE DETERMINED.
C CORRESPONDENCE
C POINTER N1 N2 N3 N4 N5 N6 N7 N8
C ARRAY VX VY VZ KFA AX AY AZ KV
C
2 ERROR=.FALSE.
READ (IN,7) MAXV,MAXS,DENS
WRITE (OT,6) HED,MAXV,MAXS,DENS
N1=1
N2=N1+MAXV
N3=N2+MAXV
N4=1
N5=N3+MAXV
N6=N5+MAXS
N7=N6+MAXS
N8=N4+MAXS
C
C VERTEX COORDINATES AND FACE LISTS ARE INPUT
C
CALL FACES ( KV, KFA, VX, VY, VZ )
CALL FACES (KA(N8),KA(N4),A(N1),A(N2),A(N3))
IF (ERROR) GO TO 3
C
C MASS PROPERTIES ARE CALCULATED
C
CALL PROPS ( KV, KFA, AX, AY, AZ,
. VX, VY, VZ )
CALL PROPS (KA(N8),KA(N4),A(N5),A(N6),A(N7),
. A(N1),A(N2),A(N3))
IF (ERROR) GO TO 3
C
C AND PRINTED OUT.
C
CALL MASSPR
GO TO 1
C

```

```

C ERROR ENCOUNTERED.  ERROR CODE PRINTED OUT. 750
C 760
  3 WRITE (OT,5) HED,NERR 770
  GO TO 1 780
C 790
C FORMAT STATEMENTS 800
C 810
  4 FORMAT (80A1) 820
  5 FORMAT (1H1,80A1///27H ERROR ENCOUNTERED, CODE = ,I1) 830
  6 FORMAT (1H1,80A1/// 7H MAXV =,I3/ 840
    . 7H MAXS =,I3/ 850
    . 7H DENS =,F10.3) 860
  7 FORMAT (2I5,F10.10) 870
  END 880
  SUBROUTINE FACES (KF,KFA,VX,VY,VZ) 890
C 900
C PURPOSE 910
C TO READ VERTEX COORDINATES AND FACE LIST DATA AND SETUP 920
C DATA IN GLOBAL STORAGE. 930
C 940
C DESCRIPTION 950
C A FACE BOUNDARY IS MADE UP OF A CLOSED SET OF NUMBERED 960
C POINTS (VERTICES) CONNECTED BY STRAIGHT LINES. 970
C A FACE LIST, ENUMERATING THE VERTICES OF THE FACE TAKEN 980
C IN SUCCESSION AROUND THE BOUNDARY, IS THE DATA 990
C REPRESENTATION OF THE FACE. 1000
C A SURFACE OF THE SOLID MAY BE COMPOSED OF MORE THAN ONE 1010
C FACE, AN INITIAL FACE AND SECONDARY FACE(S). THEY ARE 1020
C LINKED TOGETHER IN STORAGE BY A POINTER FROM ONE FACE 1030
C TO ANOTHER. 1040
C THERE ARE NO PRACTICAL RESTRICTIONS ON THE NUMBER OF 1050
C VERTICES PER FACE. THE MAXIMUM NUMBER OF SURFACES IS 1060
C CONTROLLED BY THE EXTENT OF AVAILABLE COMPUTER STORAGE 1070
C AND IS ALSO DEPENDENT ON THE TOTAL NUMBER OF VERTICES. 1080
C 1090
  COMMON  MAXS,MAXV, 1100
  . AREA,WEIGHT,VOL,DENS, 1110
  . CG(3),IO(3),ICG(3),PR(3),PRCG(3), 1120
  . ERROR,NERR, 1130
  . HED(80),IN,OT,KT, 1140
  . N8,MAXKV 1150
  INTEGER HED 1160
  DIMENSION KF(1),KFA(1),VX(1),VY(1),VZ(1) 1170
  LOGICAL ERROR 1180
  INTEGER OT 1190
C 1200
C 1210
C VERTEX COORDINATES ARE READ IN AND PRINTED. 1220
C 1230
  KT=50 1240
  1 READ (IN,14) IEND,NV,X,Y,Z 1250
  CALL COORDS (NV,X,Y,Z) 1260
  IF (NV .LE. 0 .OR. NV .GT. MAXV) GO TO 9 1270
  VX(NV)=X 1280
  VY(NV)=Y 1290
  VZ(NV)=Z 1300
  IF (IEND .EQ. 0) GO TO 1 1310
C 1320
C FACE POINTER ARRAY IS CLEARED. 1330
C 1340
  DO 2 I=1,MAXS 1350
  2 KFA(I)=0 1360
  KT=50 1370
  M1=1 1380
  3 M2=M1+2 1390
  M3=M1+16 1400
C 1410
C SURFACE ELEMENT DEFINITIONS ARE INPUT AND STORED. 1420
C 1430
  READ (IN,15) IEND,KF(M1),(KF(M),M=M2,M3) 1440
  NF=KF(M1) 1450
  K=KFA(NF) 1460
  KF(M1+1)=0 1470
  IF (K .NE. 0) GO TO 4 1480
C INITIAL (OR ONLY) FACE DEFINITION. 1490
  KFA(NF)=M1 1500

```



```

      GO TO 5
C SECONDARY FACE POINTER IS STORED.
  4  KK=K
      K=KF(KK+1)
      IF (K .NE. 0) GO TO 4
      K=KF(KK+1)
      IF (K .NE. 0) GO TO 4
      KF(KK+1)=M1
C
C CHECK IS MADE TO SEE IF DEFINITION IS COMPLETE.
C
  5  KV1=KF(M2)
      M2=M2+1
  6  DO 7 I=M2,M3
      KV=KF(I)
      IF (KV .EQ. KV1) GO TO 8
      IF (KV .LE. 0 .OR. KV .GT. MAXV) GO TO 8
  7  CONTINUE
C
C DEFINITION CONTINUED.
C
      M2=M3+1
      M3=M3+16
      READ (IN,15) IEND, (KF(M),M=M2,M3)
      GO TO 6
C
C GETTING SET FOR NEXT SURFACE, PRINT CURRENT SET OF VERTICES.
C
  8  M2=M1+2
      M3=I
      CALL FACVRT(NF,KF,M2,M3)
      IF (KV .LE. 0 .OR. KV .GT. MAXV) GO TO 11
      M1=I+1
      IF (M1+N8 .GT. MAXKV) GO TO 12
      IF (IEND .EQ. 0) GO TO 3
      RETURN
C
C ERROR RETURN.
C
C VERTEX NUMBER OUT OF RANGE
  9  NERR=1
      GO TO 13
C INPUT FACE NUMBER OUT OF RANGE
 10  NERR=2
      GO TO 13
C ILLEGAL VERTEX NUMBER.
 11  NERR=3
      GO TO 13
C STORAGE EXCEEDED.
 12  NERR=4
 13  ERROR=.TRUE.
      RETURN
C
C FORMAT STATEMENTS
C
 14  FORMAT (I1,I4,5X,3F10.0)
 15  FORMAT (I1,I4,15I5)
      END
      SUBROUTINE COORDS (NV,X,Y,Z)
C
C PURPOSE
C PRINTS VERTEX COORDINATES, WITH PAGE CONTROL.
C
C
      COMMON      MAXS,MAXV,
      .          AREA,WEIGHT,VOL,DENS,
      .          CG(3),IO(3),ICG(3),PR(3),PRCG(3),
      .          ERROR,NERR,
      .          HED(80),IN,OT,KT
      INTEGER OT
      INTEGER HED
C
C
      IF (KT .LT. 50) GO TO 1
C TITLE AND HEADING.
      WRITE (OT,2) HED

```

KT=0	2270
WRITE (OT,3)	2280
C VERTEX COORDINATES	2290
1 KT=KT+1	2300
WRITE (OT,4) NV,X,Y,Z	2310
RETURN	2320
C	2330
C FORMAT STATEMENTS.	2340
C	2350
2 FORMAT (1H1,80A1///)	2360
3 FORMAT (7H VERTEX,9X,1HX,12X,1HY,12X,1HZ /)	2370
4 FORMAT (I6,3F13.2)	2380
END	2390
SUBROUTINE FACVRT (NF,KV,M2,M3)	2400
C	2410
C PURPOSE	2420
C PRINT FACE VERTICES, WITH PAGE CONTROL.	2430
C	2440
C	2450
COMMON MAXS,MAXV,	2460
. AREA,WEIGHT,VOL,DENS,	2470
. CG(3),IO(3),ICG(3),PR(3),PRCG(3),	2480
. ERROR,NERR,	2490
. HED(80),IN,OT,KT	2500
DIMENSION KV(1)	2510
INTEGER HED	2520
INTEGER OT	2530
C	2540
C	2550
IF (KT .LT. 50) GO TO 1	2560
C TITLE AND HEADING.	2570
WRITE (OT,2) HED	2580
WRITE (OT,3)	2590
KT=0	2600
C	2610
C FACE VERTICES.	2620
C	2630
1 KT=KT+1	2640
M4=MIN0(M3,M2+9)	2650
WRITE (OT,4) NF, (KV(M),M=M2,M4)	2660
IF (M4 .EQ. M3) RETURN	2670
M4=M2+10	2680
WRITE (OT,5) (KV(M),M=M4,M3)	2690
RETURN	2700
C	2710
C FORMAT STATEMENTS.	2720
C	2730
2 FORMAT (1H1,80A1///)	2740
3 FORMAT (25H FACE VERTICES . . . /)	2750
4 FORMAT (/I5,5H - ,10I5)	2760
5 FORMAT (10X,10I5)	2770
END	2780
SUBROUTINE MASSPR	2790
C	2800
C PURPOSE	2810
C PRINT MASS PROPERTIES OF A SOLID.	2820
C	2830
C DESCRIPTION	2840
C PRINTS	2850
C DENSITY,AREA,VOLUME,WEIGHT	2860
C CENTER OF GRAVITY COORDINATES	2870
C COORDINATES OF MASS MOMENTS OF INERTIA W/R TO ORIGIN	2880
C COORDINATES OF MASS MOMENTS OF INERTIA W/R TO CG	2890
C COORDINATES OF PROD MOMENTS OF INERTIA W/R TO ORIGIN	2900
C COORDINATES OF PROD MOMENTS OF INERTIA W/R TO CG	2910
C	2920
COMMON MAXS,MAXV,	2930
. AREA,WEIGHT,VOL,DENS,	2940
. CG(3),IO(3),ICG(3),PR(3),PRCG(3),	2950
. ERROR,NERR,	2960
. HED(80),IN,OT,KT	2970
INTEGER HED	2980
INTEGER OT	2990
REAL IO,ICG	3000
C	3010
C	3020

```

WRITE (OT,1) HED 3030
WRITE (OT,2) DENS,AREA,VOL,WEIGHT 3040
WRITE (OT,3) 3050
WRITE (OT,4) CG 3060
WRITE (OT,5) IO 3070
WRITE (OT,6) ICG 3080
WRITE (OT,8) 3090
WRITE (OT,7) PR 3100
WRITE (OT,6) PRCG 3110
RETURN 3120
C 3130
C FORMAT STATEMENTS. 3140
C 3150
1 FORMAT (1H1,80A1///) 3160
2 FORMAT (10X,7HDENSITY,9X,4HAREA,10X,6HVOLUME,9X, 3170
. 3180
6HWIGHT//2X,4F15.5///)
3 FORMAT (13X,4HAXES,8X,6HX-COMP,9X,6HY-COMP,9X, 3190
. 3200
6HZ-COMP /)
4 FORMAT (1X,16H CG COORD,3F15.5/) 3210
5 FORMAT (1X,16HMOMENT COORD,3F15.5 ) 3220
6 FORMAT (1X,16H CG,3F15.5/) 3230
7 FORMAT (1X,16HPROD MOM COORD,3F15.5) 3240
8 FORMAT (//13X,4HAXES,8X,7HXY-COMP,8X,7HYZ-COMP,8X, 3250
.7HZX-COMP/) 3260
END 3270
C SOLID POLYHEDRON MEASURES 3280
C ----- 3290
C 3300
C 3310
C PURPOSE 3320
C TO CALCULATE MASS PROPERTIES OF THREE DIMENSIONAL SOLIDS 3330
C BY A SURFACE INTEGRATION TECHNIQUE USING FOUR-POINT 3340
C GAUSSIAN QUADRATURE OVER TRIANGLES. 3350
C 3360
C SOLID 3370
C DEFINED BY A SET OF SURFACES, WHERE EACH SURFACE IS 3380
C COMPOSED OF ONE OR MORE COPLANAR CLOSED POLYGONAL 3390
C PATCHES CALLED FACES. DISCONNECTED FACES AND FACES WITH 3400
C ONE OR MORE HOLES IN THEM MAY BE USED, IF NECESSARY, 3410
C TO REPRESENT THE SOLID. 3420
C 3430
C FACE 3440
C A FACE IS REPRESENTED BY THE SET OF VERTEX POINTS OF THE 3450
C POLYGONAL SURFACE PATCH, ORDERED BY STARTING WITH AN 3460
C ARBITRARY VERTEX, AND PROCEEDING POINT BY POINT AROUND 3470
C THE BOUNDARY UNTIL THE FIRST POINT IS REACHED. 3480
C 3490
C FACE LIST 3500
C THE LIST OF THE SET OF VERTEX NUMBERS, ORDERED AS ABOVE 3510
C FOR A FACE, IS CALLED A FACE LIST, WHICH IS THE DATA 3520
C REPRESENTATION OF A FACE. A FACE LIST HAS THE FORM 3530
C NSF,NA,NB,NC,...,NI,...,NA 3540
C WHERE NSF IS THE SURFACE NUMBER AND NA,NB,NC,NI ARE 3550
C VERTEX NUMBERS. NOTE THAT THE FACE LIST IS A CLOSED 3560
C LIST, STARTING AND ENDING WITH THE FIRST POINT SELECTED. 3570
C FOR ANY SURFACE WITH ONE OR MORE HOLES IN 3580
C IT, ANY INTERIOR BOUNDARY (HOLE) IS REPRESENTED BY A 3590
C FACE LIST WITH VERTICES ENUMERATED IN A SEQUENCE 3600
C OPPOSITE TO THAT GIVEN FOR AN EXTERIOR BOUNDARY. 3610
C 3620
C INITIAL, SECONDARY FACE LISTS 3630
C WHEN MULTIPLE FACE LISTS ARE NEEDED TO DEFINE A SURFACE, 3640
C (IE, A SURFACE WITH A HOLE OR ONE WITH DISCONNECTED BUT 3650
C COPLANAR PARTS), THE FIRST LIST TO APPEAR IN THE DATA IS 3660
C CALLED THE INITIAL FACE LIST, AND ANY OTHER IS CALLED A 3670
C SECONDARY FACE LIST. THESE LISTS WILL HAVE A COMMON 3680
C SURFACE NUMBER AND WILL BE LINKED TOGETHER IN STORAGE BY 3690
C THE DATA ENTRY ROUTINE. 3700
C 3710
C GLOBAL DATA ARRAYS 3720
C FACE LIST DATA IS STORED COMPACTLY IN A LARGE OPEN-ENDED 3730
C DATA ARRAY KV. 3740
C A LINKAGE ARRAY, KFA, SETUP BY THE DATA ENTRY ROUTINE, 3750
C CONTAINS POINTERS TO THE INITIAL FACE LIST (IN KV) FOR 3760
C EACH SURFACE, ORDERED BY SURFACE NUMBER. IF L1, L2, L3 3770
C ARE LOCATIONS IN KV OF FACE LISTS FOR SURFACES 1, 2, 3, 3780

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```

C THEN THE PROGRAM WOULD SET 3790
C   KFA(1)=L1 3800
C   KFA(2)=L2 3810
C   KFA(3)=L3 3820
C A POINTER VALUE OF ZERO (KFA(I)=0) INDICATES SURFACE I 3830
C WAS NOT UTILIZED IN DEFINING THE SOLID. HENCE GAPS MAY 3840
C OCCUR IN THE SURFACE NUMBERS. ALSO, THE ORDER OF DATA 3850
C ENTRY OF THE FACE LISTS CAN BE ENTIRELY ARBITRARY, 3860
C BECAUSE OF THE AFORE-MENTIONED LINKAGE MECHANISM. 3870
C 3880
C LINKAGE TO SECONDARY FACE LISTS 3890
C A FACE LIST IN THE ARRAY KV HAS THE FORM 3900
C   NSF,LSI,NA,NB,NC,...,NL,NJ,NK,...,NA 3910
C WHERE NSF IC THE SURFACE NUMBER, NA,NB,... ARE VERTEX 3920
C NUMBERS, AND LSI IS A LINK TO A SECONDARY FACE LIST. 3930
C LSI=0 INDICATES THIS IS THE LAST FACE LIST FOR A SURFACE. 3940
C NOTE, IT MAY BE THE ONLY LIST. 3950
C 3960
C INTEGRATION TECHNIQUE 3970
C THE ALGORITHM IS PERFORMED BY TWO SUBROUTINES, PROPS AND 3980
C SRFINT. 3990
C SURFACES ARE PROCESSED ONE AT A TIME AND THE MASS 4000
C PROPERTY CALCULATIONS ARE ACCUMULATED DURING THE 4010
C PROCESSING. 4020
C SURFACE PROCESSING INVOLVES STARTING WITH THE INITIAL 4030
C FACE, AND PROCEEDING THROUGH THE SECONDARY FACES, IF ANY. 4040
C SURFACE AREA VECTOR COMPONENTS ARE SUMMED FOR EACH FACE 4050
C OF THE SURFACE. THE AREA VECTOR, WHOSE MAGNITUDE IS THE 4060
C AREA OF THE SURFACE, IS THEN COMPUTED FROM THE 4070
C COMPONENTS AND IS ACCUMULATED IN THE TOTAL SURFACE 4080
C AREA OF THE SOLID. 4090
C IN THE FACE PROCESSING, FACES ARE DIVIDED INTO ALL 4100
C UNIQUE TRIANGLES HAVING AS VERTICES THE THE FIRST ONE OF 4110
C THE LIST AND ANY OTHER TWO ADJACENT VERTICES OF THE FACE 4120
C POLYGON. FOUR-POINT GAUSSIAN QUADRATURE IS APPLIED TO 4130
C EACH TRIANGLE, AND THE RESULTS ARE ACCUMULATED IN THE 4140
C TOTAL MASS PROPERTIES. THUS EACH TRIANGLE OF EACH FACE 4150
C OF EACH SURFACE HAS A CONTRIBUTION IN THE FINAL RESULTS. 4160
C 4170
C GLOBAL NOMENCLATURE 4180
C AREA = SURFACE AREA OF SOLID 4190
C AX,AY, 4200
C   AZ = SURFACE AREA VECTOR COMPONENTS ARRAYS 4210
C CG = CENTER OF GRAVITY COORDINATE ARRAY 4220
C DENS = MASS DENSITY OF SOLID 4230
C ERROR = LOGICAL ERROR INDICATOR 4240
C ICG = COORDS OF MASS MOMENTS, WITH RESPECT TO CG 4250
C IO = COORDS OF MASS MOMENTS, WITH RESPECT TO ORIGIN 4260
C KFA = POINTER ARRAY. KFA(I) POINTS TO LOCATION IN KV 4270
C   OF INITIAL FACE LIST FOR SURFACE I. 4280
C KV = FACE LISTS STORAGE ARRAY 4290
C MAXS = NUMBER OF SURFACES ON SOLID 4300
C MAXV = NUMBER OF VERTEX POINTS USED IN DEFINING SOLID 4310
C NERR = ERROR CODE 4320
C NS = CURRENT SURFACE NUMBER 4330
C PR = COORDS OF PROD MOMENTS, WITH RESPECT TO ORIGIN 4340
C PRCG = COORDS OF PROD MOMENTS, WITH RESPECT TO CG 4350
C VOL = VOLUME OF SOLID 4360
C VX,VY, 4370
C   VZ = VERTEX COORDINATE ARRAYS 4380
C WEIGHT = WEIGHT OF SOLID 4390
C 4400
C   SUBROUTINE PROPS (KV,KFA,AX,AY,AZ,VX,VY,VZ) 4410
C 4420
C PURPOSE 4430
C CALCULATES MASS PROPERTIES OF A GENERAL THREE 4440
C DIMENSIONAL SOLID USING FOUR POINT GAUSSIAN QUADRATURE 4450
C OVER TRIANGLES. 4460
C 4470
C DESCRIPTION 4480
C CALCULATES CONTRIBUTION OF EACH SURFACE TO MASS 4490
C PROPERTIES INTEGRALS. INTEGRATION IS PERFORMED IN 4500
C SUBROUTINE SRFINT. MASS PROPERTIES CALCULATIONS ARE 4510
C THEN COMPLETED IN THIS ROUTINE. DATA IS ASSUMED TO HAVE 4520
C BEEN PREVIOUSLY SETUP IN STORAGE BY A DATA ENTRY ROUTINE. 4530
C 4540

```

```

C LOCAL NOMENCLATURE 4550
C NS = CURRENT SURFACE BEING PROCESSED 4560
C LA = POINTER TO FACE LIST (INITIAL OR SECONDARY) DATA 4570
C 4580
COMMON MAXS,MAXV, 4590
. AREA,WEIGHT,VOL,DENS, 4600
. CG(3),IO(3),ICG(3),PR(3),PRCG(3), 4610
. ERROR,NERR 4620
DIMENSION KV(1),KFA(1),VX(1),VY(1),VZ(1) 4630
DIMENSION AX(1),AY(1),AZ(1) 4640
LOGICAL ERROR 4650
REAL IO,ICG 4660
C 4670
C CLEARING OF ACCUMULATION PARAMETERS. 4680
C 4690
AREA=0. 4700
VOL=0. 4710
DO 1 J=1,MAXS 4720
AX(J)=0. 4730
AY(J)=0. 4740
1 AZ(J)=0. 4750
DO 2 J=1,3 4760
CG(J)=0. 4770
IO(J)=0. 4780
2 PR(J)=0. 4790
C 4800
C EACH SURFACE IS PROCESSED. 4810
C 4820
DO 4 NS=1,MAXS 4830
C LA IS A POINTER TO THE FACE LIST DATA 4840
LA=KFA(NS) 4850
IF (LA .EQ. 0) GO TO 4 4860
C CONTRIBUTION OF EACH FACE OF CURRENT SURFACE IS CALCULATED 4870
3 IF (KV(LA) .NE. NS) GO TO 5 4880
CALL SRFINT (KV(LA+2),AX,AY,AZ,VX,VY,VZ,NS) 4890
C CHECK IS MADE FOR SECONDARY FACE LIST. 4900
LA=KV(LA+1) 4910
IF (LA .NE. 0) GO TO 3 4920
C AREA SUMMATION. 4930
AREA=AREA+SQRT (AX(NS)**2+AY(NS)**2+AZ(NS)**2) 4940
4 CONTINUE 4950
C 4960
C MASS PROPERTY CALCULATIONS ARE COMPLETED. 4970
C 4980
VOL=VOL/3.0 4990
WEIGHT=VOL*DENS 5000
VOL2=VOL*2. 5010
VOL3=VOL/3.0 5020
DENS2=DENS/2.0 5030
DENS3=DENS/3.0 5040
CG(1)=CG(1)/VOL2 5050
CG(2)=CG(2)/VOL2 5060
CG(3)=CG(3)/VOL2 5070
IO(1)=IO(1)*DENS3 5080
IO(2)=IO(2)*DENS3 5090
IO(3)=IO(3)*DENS3 5100
ICG(1)=IO(1)-(CG(2)**2+CG(3)**2)*WEIGHT 5110
ICG(2)=IO(2)-(CG(1)**2+CG(3)**2)*WEIGHT 5120
ICG(3)=IO(3)-(CG(1)**2+CG(2)**2)*WEIGHT 5130
PR(1)=PR(1)*DENS2 5140
PR(2)=PR(2)*DENS2 5150
PR(3)=PR(3)*DENS2 5160
PRCG(1)=PR(1)-CG(1)*CG(2)*WEIGHT 5170
PRCG(2)=PR(2)-CG(2)*CG(3)*WEIGHT 5180
PRCG(3)=PR(3)-CG(3)*CG(1)*WEIGHT 5190
RETURN 5200
C 5210
C ERROR RETURN, DATA REDUNDANCY CHECK FAILED. 5220
C 5230
5 ERROR=.TRUE. 5240
NERR=5 5250
RETURN 5260
END 5270
SUBROUTINE SRFINT (KV,AX,AY,AZ,VX,VY,VZ,NS) 5280
C 5290
C PURPOSE 5300

```

```

C TO CALCULATE SURFACE INTEGRALS OVER A FACE OF A SOLID. 5310
C 5320
C DESCRIPTION 5330
C DIVIDES FACE INTO TRIANGLES HAVING AS VERTICES THE 5340
C FIRST VERTEX OF THE POLYGON AND TWO ADJACENT VERTICES 5350
C OF THE POLYGON. NUMERICAL INTEGRATION IS PERFORMED 5360
C USING FOUR-POINT GAUSSIAN QUADRATURE OVER THE TRIANGLES. 5370
C RESULTS ARE ACCUMULATED TO GET TOTALS. 5380
C 5390
C LOCAL NOMENCLATURE 5400
C ANX, 5410
C ANY, 5420
C ANZ = AREA VECTOR COMPONENTS FOR SINGLE TRIANGLE 5430
C CX,CY, 5440
C CZ = 1ST,2ND OR 3RD POWERS OF QUADRATURE COORDS, 5450
C WEIGHTED AND AREA NORMALIZED 5460
C PX,PY, 5470
C PZ = QUADRATURE POINTS COORDINATE ARRAYS 5480
C W = WEIGHTING FACTORS ARRAY 5490
C X,Y,Z = FACE TRIANGLE COORDINATE ARRAYS 5500
C 5510
C COMMON MAXS,MAXV, 5520
C AREA,WEIGHT,VOL,DENS, 5530
C CG(3),IO(3),ICG(3),PR(3),PRCG(3) 5540
C DIMENSION KV(1),AX(1),AY(1),AZ(1),VX(1),VY(1),VZ(1) 5550
C DIMENSION X(3),Y(3),Z(3),PX(4),PY(4),PZ(4),W(4) 5560
C REAL IO,ICG 5570
C 5580
C W(I) ARE WEIGHING FACTORS FOR 4-POINT GAUSSIAN QUADRATURE 5590
C OVER TRIANGLES. 5600
C 5610
C W(1)=-.5625 5620
C W(2)=.5208333333333333 5630
C W(3)=.5208333333333333 5640
C W(4)=.5208333333333333 5650
C 5660
C THE POLYGON IS SEGMENTED INTO TRIANGLES, EACH HAVING THE 5670
C FIRST VERTEX OF THE POLYGON AS THE FIRST POINT OF THE 5680
C TRIANGLE. 5690
C 5700
C K=KV(1) 5710
C X(1)=VX(K) 5720
C Y(1)=VY(K) 5730
C Z(1)=VZ(K) 5740
C DO 3 L=2,1000 5750
C K=KV(L+1) 5760
C TESTING FOR END OF DEFINITION. 5770
C IF (K .EQ. KV(1)) GO TO 4 5780
C X(3)=VX(K) 5790
C Y(3)=VY(K) 5800
C Z(3)=VZ(K) 5810
C K=KV(L) 5820
C X(2)=VX(K) 5830
C Y(2)=VY(K) 5840
C Z(2)=VZ(K) 5850
C 5860
C DATA IS NOW IN TERMS OF A SINGLE TRIANGLE. 5870
C 5880
C PX(1)=(X(1)+X(2)+X(3))/3. 5890
C PY(1)=(Y(1)+Y(2)+Y(3))/3. 5900
C PZ(1)=(Z(1)+Z(2)+Z(3))/3. 5910
C DX=.6*PX(1) 5920
C DY=.6*PY(1) 5930
C DZ=.6*PZ(1) 5940
C DO 1 I=1,3 5950
C K=I+1 5960
C PX(K)=DX+.4*X(I) 5970
C PY(K)=DY+.4*Y(I) 5980
C PZ(K)=DZ+.4*Z(I) 5990
C 1 CONTINUE 6000
C 6010
C AREA VECTORS CALCULATED. 6020
C 6030
C ANX=.5*(Z(1)*(Y(2)-Y(3))+Z(2)*(Y(3)-Y(1)) 6040
C +Z(3)*(Y(1)-Y(2))) 6050
C ANY=.5*(X(1)*(Z(2)-Z(3))+X(2)*(Z(3)-Z(1)) 6060

```

+X(3)*(Z(1)-Z(2))	6070
ANZ=.5*(Y(1)*(X(2)-X(3))+Y(2)*(X(3)-X(1))	6080
+Y(3)*(X(1)-X(2)))	6090
AX(NS)=AX(NS)+ANZ	6100
AY(NS)=AY(NS)+ANZ	6110
AZ(NS)=AZ(NS)+ANZ	6120
C	6130
C CALCULATION OF MASS PROPERTIES	6140
C	6150
DO 2 I=1,4	6160
C 1ST POWER OF INTEGRATION COORDINATES	6170
CX=ANX*PX(I)*W(I)	6180
CY=ANY*PY(I)*W(I)	6190
CZ=ANZ*PZ(I)*W(I)	6200
VOL=VOL+(CX+CY+CZ)	6210
C 2ND POWER OF INTEGRATION COORDINATES	6220
CX=CX*PX(I)	6230
CY=CY*PY(I)	6240
CZ=CZ*PZ(I)	6250
CG(1)=CG(1)+CX	6260
CG(2)=CG(2)+CY	6270
CG(3)=CG(3)+CZ	6280
PR(1)=PR(1)+CX*PY(I)	6290
PR(2)=PR(2)+CY*PZ(I)	6300
PR(3)=PR(3)+CZ*PX(I)	6310
C 3RD POWER OF INTEGRATION COORDINATES	6320
CX=CX*PX(I)	6330
CY=CY*PY(I)	6340
CZ=CZ*PZ(I)	6350
IO(1)=IO(1)+CY+CZ	6360
IO(2)=IO(2)+CZ+CX	6370
IO(3)=IO(3)+CX+CY	6380
2 CONTINUE	6390
3 CONTINUE	6400
C	6410
4 CONTINUE	6420
RETURN	6430
END	6440

ALGORITHM 551

A Fortran Subroutine for the L_1 Solution of Overdetermined Systems of Linear Equations [F4]

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Key Words and Phrases: overdetermined system of linear equations, discrete linear L_1 approximation, linear programming, dual simplex algorithm, triangular decomposition

CR Categories: 5.13, 5.41

Language: Fortran

DESCRIPTION

This algorithm is a Fortran implementation for the procedure developed in [1] and [2]. It solves an overdetermined system of linear equations of the form

$$Ca = f,$$

in the L_1 norm. C is a given real n -by- m matrix of rank $k \leq m \leq n$ and f is a given real n -vector.

The L_1 solution to this system is the m -vector a^* which minimizes the L_1 norm

$$z = \sum_{i=1}^n |r_i|,$$

where r_i is the i th residual and is given by

$$r_i = \sum_{j=1}^m c_{ij}a_j - f_i.$$

This subroutine uses a dual simplex method and a suitable triangular decomposition method to the basis matrix. In [2] the testing of this subroutine, as well as numerical results and comments, are given.

REFERENCES

1. ABDELMALEK, N.N. An efficient method for the discrete linear L_1 approximation problem. *Math. Comput.* 29 (1975), 844-850.
2. ABDELMALEK, N.N. L_1 solution of overdetermined systems of linear equations. *ACM Trans. Math. Softw.* 6, 2 (June 1980), 220-227.

ALGORITHM

```
C      THIS PROGRAM IS A DRIVER FOR THE SUBROUTINE L1
C      WHICH SOLVES AN OVERDETERMINED SYSTEM OF LINEAR EQUATIONS
C      IN THE L1 NORM, USING A DUAL SIMPLEX METHOD.
C      THE OVERDETERMINED SYSTEM HAS THE FORM CA=F
C      C IS A GIVEN REAL N BY M MATRIX OF RANK K, K.LE.M.LE.N
C      F IS A GIVEN REAL N-VECTOR.
```

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```

C   A IS THE SOLUTION M-VECTOR.
C
C
C
      DIMENSION AA(252,27)
      DIMENSION FAY(8,250),EF(250)
      DIMENSION CT(25,250),F(250),R(250),A(25),BINV(25,25)
      DIMENSION P(350),IC(250),IB(250),Y(250),TH(250),IR(25)
      DIMENSION U(25),V(25),W(25),GINV(25,25),VB(25)
      DIMENSION BV(25),IBOUND(250),ICBAS(25),IRBAS(25),ZC(250)
      DIMENSION F1(50),F2(50),BA(6,15)
      INTEGER S(250)
C
101  FORMAT(10H1EXAMPLE ,I3,8H SIZE ,I3,8H BY ,I3)
102  FORMAT(40H0 R.H.S. OF THE SYSTEM )
103  FORMAT(40H0TRANSPOSE OF COEFFICIENT MATRIX )
104  FORMAT(17H0EXECUTION TIME =,F12.5,13HSECONDS )
105  FORMAT(9H0L1 NORM=,F10.5,6H,RANK=,I4,6H,ITER=,I4,5H,IND=,I4)
106  FORMAT(40H0THE ANSWER : A(I) OR X(I) )
107  FORMAT(40H0RESIDUALS R(J) OR E(J) )
108  FORMAT(40H0IC(I) )
109  FORMAT(40H0IR(I) )
110  FORMAT(1H ,10F12.5)
111  FORMAT(1H ,20I4)
115  FORMAT(9H0L1 NORM=,F10.5,6H,RANK=,F3.0,6H,ITER=,F3.0,5H,IND=,F2.0)
120  FORMAT(1H ,8F15.5)
125  FORMAT(5F10.6)
126  FORMAT(6F10.6)
C
      CALL SUNDER
      PREC=1.E-6
      EPS=1.E-4
      TOLER=1.0E-4
      MM=25
      MM2=MM+2
      MMM=(MM*(MM+3))/2
      NN=250
      NN2=NN+2
      IEXMPL=0
1  IEXMPL=IEXMPL+1
      GO TO (2,4,5,6,7,8,9,13,17,19,20,21,22,26,27,28,29,32,35,100),
      *IEXMPL
C EXAMPLE 1.
2  N=201
      DX=0.02
      DO 3 I=1,N
      X1=DX*FLOAT(I-1)
      X2=X1*X1
      X3=X2*X1
      X4=X2*X2
      X5=X3*X2
      X6=X3*X3
      X7=X3*X4
      FAY(1,I)=1.
      FAY(2,I)=X1
      FAY(3,I)=X2
      FAY(4,I)=X3
      FAY(5,I)=X4
      FAY(6,I)=X5
      FAY(7,I)=X6
3  EF(I)= (EXP(-X1))*SIN(X1)
      M=1
      GO TO 10
4  M=2
      GO TO 10
5  M=3
      GO TO 10
6  M=4
      GO TO 10
7  M=5
      GO TO 10
8  M=6
      GO TO 10
9  M=7
10 DO 12 J=1,N
      F(J)=EF(J)

```

```

      DO 11 I=1,M
11    CT(I,J)=FAY(I,J)
12  CONTINUE
      GO TO 90
C EXAMPLE 2.
13  N=100
      N2=50
      M=6
      READ(1,125) (F1(I), I=1,N2)
      READ(1,125) (F2(I), I=1,N2)
      DO 14 J=1,N2
      F(J)=F1(J)
      JN2=J+N2
14  F(JN2)=F2(J)
      DO 16 J=1,100
      D=0.1*FLOAT(J)
      DD=D*D
      DO 15 I=1,100
      E=0.1*FLOAT(I)
      EE=E*E
      K=10*(J-1)+I
      CT(1,K)=1.0
      CT(2,K)=E
      CT(3,K)=D
      CT(4,K)=E*D
      CT(5,K)=EE
      CT(6,K)=DD
15  CONTINUE
16  CONTINUE
      GO TO 90
C EXAMPLE 3.
17  N=51
      M=8
      DO 18 I=1,N
      D=0.02*FLOAT(I-1)
      DD=D*D
      DDD=D*DD
      E1=D-0.1
      E2=D-0.2
      E3=D-0.4
      E4=D-0.7
      E13=E1*E1*E1
      E23=E2*E2*E2
      E33=E3*E3*E3
      E43=E4*E4*E4
      CT(1,I)=1.0
      CT(2,I)=D
      CT(3,I)=DD
      CT(4,I)=DDD
      IF(E1.GT.0.0) CT(5,I)=E13
      IF(E1.LE.0.0) CT(5,I)=0.
      IF(E2.GT.0.0) CT(6,I)=E23
      IF(E2.LE.0.0) CT(6,I)=0.
      IF(E3.GT.0.0) CT(7,I)=E33
      IF(E3.LE.0.0) CT(7,I)=0.
      IF(E4.GT.0.0) CT(8,I)=E43
      IF(E4.LE.0.0) CT(8,I)=0.
18  F(I)= SQRT(D)
      GO TO 90
C EXAMPLE 4.
19  M=6
      PI=4.0*ATAN(1.0)
      X01=PI/2.0
      X02=X01*X01
      X03=X02*X01
      X04=X02*X02
      N=23
      GO TO 23
20  N=53
      GO TO 23
21  N=103
      GO TO 23
22  N=203
23  N2=N-2
      N3=N-3
      DX=PI/FLOAT(N3)

```

```

DO 24 I=1,N2
X1=-X01+DX*FLOAT(I-1)
X2=X1*X1
X3=X2*X1
X4=X2*X2
X5=X3*X2
CT(1,I)=1.0
CT(2,I)=X1
CT(3,I)=X2
CT(4,I)=X3
CT(5,I)=X4
CT(6,I)=X5
24 F(I)= SIN(X1)
G=-1.0
DO 25 I=1,2
IF(I.EQ.2) G=1.0
J=N2+I
CT(1,J)=0.0
CT(2,J)=1000.0
CT(3,J)=2000.0*C*X01
CT(4,J)=3000.0*X02
CT(5,J)=4000.0*C*X03
CT(6,J)=5000.0*X04
25 F(J)=0.0
GO TO 90
C EXAMPLE 5.
26 M=11
D=20.0
DX=1.0/D
N=21
GO TO 30
27 D=50.0
DX=1.0/D
N=51
GO TO 30
28 D=100.0
DX=1.0/D
N=101
GO TO 30
29 D=200.0
DX=1.0/D
N=201
30 DO 31 I=1,N
X1=DX*FLOAT(I-1)
X2=X1+X1
X3=X1+X2
X4=X2+X2
X5=X2+X3
CT(1,I)=1.0
CT(2,I)= SIN(X1)
CT(3,I)= COS(X1)
CT(4,I)= SIN(X2)
CT(5,I)= COS(X2)
CT(6,I)= SIN(X3)
CT(7,I)= COS(X3)
CT(8,I)= SIN(X4)
CT(9,I)= COS(X4)
CT(10,I)= SIN(X5)
CT(11,I)= COS(X5)
G= EXP(X1)
G1= EXP(0.5)
F(I)=G1
IF(G.LT.G1) F(I)=G
31 CONTINUE
GO TO 90
C EXAMPLE 6.
32 M=5
M1=M+1
N=15
DO 34 J=1,N
READ(1,126) (BA(I,J),I=1,M1)
DO 33 I=1,M
33 CT(I,J)=BA(I,J)
F(J)=BA(M1,J)
34 CONTINUE
GO TO 90

```

```

C EXAMPLE 7
  35 M=5
      N=51
      DX=0.02
      DO 36 I=1,N
      X1=DX*FLOAT(I-1)
      X2=X1*X1
      X3=X2*X1
      X4=X2*X2
      CT(1,I)=1.0
      CT(2,I)=X1
      CT(3,I)=X2
      CT(4,I)=X3
      CT(5,I)=X4
      GX=0.0
      IF(X1.GT.0.93) GX=5.0
      F(I)=X1*(1.0+X1*(1.0+X1*(1.0+X1))) +GX
  36 CONTINUE

C
  90 WRITE (3,101) IEXMPL,N,M
C   WRITE(3,102)
C   WRITE(3,110) (F(J),J=1,N)
C   WRITE(3,103)
C   DO 93 I=1,M
C  93 WRITE(3,110) (CT(I,J),J=1,N)
      CALL SETCLK
      CALL L1(M,N,MM,MMM,NN,CT,F,PREC,EPS,IC,IR,IB,
      *U,V,W,Y,TH,P,GINV,VB,IRANK,ITER,R,A,Z,IND)
      CALL RDCLK(TX)
      WRITE(3,108)
      WRITE(3,111) (IC(I),I=1,M)
      WRITE(3,109)
      WRITE(3,111) (IR(I),I=1,M)
      WRITE(3,107)
      WRITE(3,110) (R(J),J=1,N)
      WRITE(3,106)
      WRITE(3,120) (A(I),I=1,M)
      WRITE(3,105) Z,IRANK,ITER,IND
      WRITE(3,104) TX
      GO TO 1
  100 STOP
      END

      SUBROUTINE PSLV(ID,K,MM,MMM,P,B,X)
C   THIS SUBROUTINE SOLVES THE SQUARE NON-SINGULAR SYSTEM
C OF LINEAR EQUATIONS
C
C
C   P*X=B,
C OR THE SQUARE NON-SINGULAR SYSTEM OF LINEAR EQUATIONS
C
C   P(TRANPOSE)*X=B,
C WHERE P IS AN UPPER TRIANGULAR MATRIX, B IS THE RIGHT HAND
C SIDE VECTOR AND X IS THE SOLUTION VECTOR.
C
C   THE INPUT DATA TO THE SUBROUTINE.
C ID   AN INTEGER INDICATOR SPECIFIED BY THE USER.
C     IF ID=1 THE EQUATION P*X=B IS SOLVED.
C     IF ID= ANY INTEGER OTHER THAN 1, THE EQUATION
C     P(TRANPOSE)*X=B IS SOLVED.
C K    AN INTEGER = THE NUMBER OF EQUATIONS OF THE GIVEN
C     SYSTEM.
C MM   AN INTEGER GREATER THAN OR EQUAL TO K.
C MMM  AN INTEGER = (MM*(MM+3))/2
C P    AN MMM-VECTOR. ITS FIRST (K+1) ELEMENTS CONTAIN THE
C     FIRST K ELEMENTS OF ROW 1 OF THE UPPER TRIANGULAR
C     MATRIX PLUS AN EXTRA LOCATION TO THE RIGHT. ITS
C     NEXT K ELEMENTS CONTAIN THE (K-1) ELEMENTS OF ROW 2
C     OF THE UPPER TRIANGULAR MATRIX PLUS AN EXTRA
C     LOCATION TO THE RIGHT, ..., ETC.
C B    AN MM-VECTOR. ITS FIRST K ELEMENTS CONTAIN THE
C     R.H.S. VECTOR OF THE GIVEN SYSTEM.
C
C     THE OUTPUT OF THE SUBROUTINE.
C X    AN MM-VECTOR. ON EXIT, ITS FIRST K ELEMENTS CONTAIN
C     THE SOLUTION TO THE GIVEN SYSTEM.
C
C     DOUBLE PRECISION S,SA,SB
C     DIMENSION P(MMM),B(MM),X(MM)
C     IF(ID.NE.1) GO TO 3
C SOLUTION OF THE UPPER TRIANGULAR SYSTEM.

```

```

PSLV0010
PSLV0020
PSLV0030
PSLV0040
PSLV0050
PSLV0060
PSLV0070
PSLV0080
PSLV0090
PSLV0100
PSLV0110
PSLV0120
PSLV0130
PSLV0140
PSLV0150
PSLV0160
PSLV0170
PSLV0180
PSLV0190
PSLV0200
PSLV0210
PSLV0220
PSLV0230
PSLV0240
PSLV0250
PSLV0260
PSLV0270
PSLV0280
PSLV0290
PSLV0300
PSLV0310
PSLV0320

```

```

L=(K-1)+(K*(K+1))/2
X(K)=B(K)/P(L)
IF(K.EQ.1) RETURN
KD=3
KM1=K-1
DO 2 I=1,KM1
  J=K-I
  L=L-KD
  KD=KD+1
  S=B(J)
  LL=L
  JJ=J
  DO 1 KK=1,I
    LL=LL+1
    SA=P(LL)
    JJ=JJ+1
    SB=X(JJ)
1    S=S-SA*SB
    X(J)=S/P(L)
2  CONTINUE
RETURN
C SOLUTION OF THE LOWER TRIANGULAR SYSTEM.
3 X(1)=B(1)/P(1)
IF(K.EQ.1) RETURN
L=1
KD=K+1
DO 5 I=2,K
  L=L+KD
  KD=KD-1
  S=B(I)
  KK=I
  KKM1=I-1
  KKD=K
  DO 4 J=1,KKM1
    SA=P(KK)
    SB=X(J)
    S=S-SA*SB
    KK=KK+KKD
    KKD=KKD-1
4  CONTINUE
5  X(I)=S/P(L)
RETURN
END

```

```

PSLV0330
PSLV0340
PSLV0350
PSLV036
PSLV0370
PSLV0380
PSLV0390
PSLV0400
PSLV0410
PSLV0420
PSLV0430
PSLV0440
PSLV0450
PSLV0460
PSLV0470
PSLV0480
PSLV0490
PSLV0500
PSLV0510
PSLV0520
PSLV0530
PSLV0540
PSLV0550
PSLV0560
PSLV0570
PSLV0580
PSLV0590
PSLV0600
PSLV0610
PSLV0620
PSLV0630
PSLV0640
PSLV0650
PSLV0660
PSLV0670
PSLV0680
PSLV0690
PSLV0700
PSLV0710
PSLV0720
PSLV0730
PSLV0740
PSLV0750

```

```

SUBROUTINE L1(M,N,MM,MM,MM,NN,CT,F,PREC,EPS,IC,IR,IB,
*UF,BP,XP,T,ALFA,P,GINV,VB,IRANK,ITER,R,A,Z,IND)
C THIS SUBROUTINE SOLVES AN OVERDETERMINED SYSTEM OF
C LINEAR EQUATIONS IN THE L1 NORM BY USING A DUAL SIMPLEX
C ALGORITHM TO THE LINEAR PROGRAMMING FORMULATION OF THE
C GIVEN PROBLEM. IN THIS ALGORITHM CERTAIN INTERMEDIATE
C SIMPLEX ITERATIONS ARE SKIPPED.
C FOR PURPOSE OF NUMERICAL STABILITY, THIS SUBROUTINE
C USES A TRIANGULAR DECOMPOSITION TO THE BASIS MATRIX.
C THE SYSTEM OF LINEAR EQUATIONS HAS THE FORM
C C*A=F,
C WHERE C IS A GIVEN REAL N BY M MATRIX OF RANK K.LE.M.LE.N
C AND F IS A GIVEN REAL N-VECTOR.
C THE PROBLEM TO BE SOLVED IS TO CALCULATE THE ELEMENTS
C OF THE M-VECTOR A* WHICH GIVES THE MINIMUM NORM Z.
C Z= ABS(R(1)) + ABS(R(2)) + ... + ABS(R(N)) ,
C WHERE R(I) IS THE ITH RESIDUAL AND IS GIVEN BY
C R(I)=C(I,1)*A(1)+C(I,2)*A(2)+ ... +C(I,M)*A(M)-F(I).
C SUBROUTINE L1 IS COMPLETELY SELF CONTAINED (CONSISTS
C OF TWO SUBROUTINES L1 AND PSLV).
C THE INPUT DATA TO THE SUBROUTINE.
C M THE NUMBER OF COLUMNS OF MATRIX C.
C N THE NUMBER OF ROWS OF MATRIX C.
C MM AN INTEGER GREATER THAN OR EQUAL TO M.
C MMM AN INTEGER = (MM*(MM+3))/2 .
C NN AN INTEGER GREATER THAN OR EQUAL TO N.
C CT A MATRIX OF DIMENSIONS MM BY NN. ON ENTRY, ITS FIRST
C M ROWS AND FIRST N COLUMNS CONTAIN THE TRANSPOSE OF
C MATRIX C IN THE GIVEN SYSTEM C*A=F.
C F AN NN-VECTOR. ON ENTRY, ITS FIRST N ELEMENTS CONTAIN
C THE R.H.S. OF THE GIVEN SYSTEM C*A=F.

```

```

L1 0010
L1 0020
L1 0030
L1 0040
L1 0050
L1 0060
L1 0070
L1 0080
L1 0090
L1 0100
L1 0110
L1 0120
L1 0130
L1 0140
L1 0150
L1 0160
L1 0170
L1 0180
L1 0190
L1 0200
L1 0210
L1 0220
L1 0230
L1 0240
L1 0250
L1 0260
L1 0270
L1 0280
L1 0290
L1 0300
L1 0310

```

```

C PREC THE ROUND-OFF LEVEL OF THE COMPUTER. FOR THE IBM L1 0320
C 360/67 COMPUTER, PREC IS ABOUT 1.E-6 AND 1.E-16 FOR L1 0330
C SINGLE AND DOUBLE PRECISION RESPECTIVELY. L1 0340
C EPS A SPECIFIED TOLERANCE SUCH THAT A CALCULATED NUMBER L1 0350
C X IS CONSIDERED ZERO IF ABS(X) < EPS. FOR THE IBM 360/67 L1 0360
C COMPUTER, EPS IS USUALLY TAKEN 1.E-4 AND 1.E-11 L1 0370
C RESPECTIVELY. L1 0380
C THE RESULTS OF THE PROBLEM. L1 0390
C IRANK THE CALCULATED RANK OF MATRIX C. L1 0400
C ITER THE NUMBER OF ITERATIONS WHICH THE SOLUTION NEEDED. L1 0410
C A AN MM-VECTOR. ITS FIRST M ELEMENTS ARE THE SOLUTION L1 0420
C VECTOR A*. L1 0430
C R AN NN-VECTOR. ITS FIRST N ELEMENTS CONTAIN THE L1 0440
C THE RESIDUAL VECTOR R=C*A-F. L1 0450
C Z THE MINIMUM L1 NORM OF THE RESIDUAL VECTOR R. L1 0460
C IND A RETURN INDICATOR. IND=0 INDICATES THAT THE L1 0470
C SOLUTION VECTOR A* IS UNIQUE. IND=1, INDICATES THAT L1 0480
C A* IS MOST PROBABLY NOT UNIQUE. IND=-1 INDICATES L1 0490
C PREMATURE TERMINATION OF THE CALCULATION DUE TO L1 0500
C VERY ILL-CONDITIONING OF MATRIX C. L1 0510
C THE MEANING OF THE OTHER PARAMETERS. L1 0520
C GINV AN MM-SQUARE MATRIX. ITS FIRST IRANK COLUMNS AND L1 0530
C FIRST IRANK ROWS CONTAIN THE INVERSE OF THE INITIAL L1 0540
C BASIS MATRIX AND ITS UPDATE. L1 0550
C VB AN MM-VECTOR. ITS FIRST IRANK ELEMENTS CONTAIN THE L1 0560
C INITIAL BASIC SOLUTION AND ITS UPDATE. L1 0570
C T AN NN-VECTOR. ITS FIRST N ELEMENTS CONTAIN THE L1 0580
C ELEMENTS OF THE ROW IN THE SIMPLEX TABLEAU, THAT L1 0590
C CORRESPONDS TO THE COLUMN WHICH LEAVES THE BASIS. L1 0600
C ALFA AN NN-VECTOR. ITS FIRST N ELEMENTS CONTAIN THE L1 0610
C RATIOS: ALFA(J) = R(J)/T(J). L1 0620
C IC AN NN-VECTOR. ITS FIRST N ELEMENTS CONTAIN THE L1 0630
C COLUMN INDICES OF MATRIX CT. L1 0640
C IR AN MM VECTOR. ITS FIRST IRANK ELEMENTS CONTAIN THE L1 0650
C ROW INDICES OF THE LINEARLY INDEPENDENT ROWS OF CT. L1 0660
C IB A SIGN NN-VECTOR. ITS FIRST N ELEMENTS HAVE THE L1 0670
C VALUES +1 OR -1. IB(J)=+1 INDICATES THAT COLUMN J L1 0680
C IN THE TABLEAU IS AT ITS LOWER BOUND 0. IB(J)=-1 L1 0690
C INDICATES THAT COLUMN J IS AT ITS UPPER BOUND 2. L1 0700
C P AN MMM-VECTOR. ITS FIRST ((IRANK*(IRANK+3))/2)-1 L1 0710
C ELEMENTS CONTAIN THE (IRANK*(IRANK+1))/2 ELEMENTS OF L1 0720
C THE UPPER TRIANGULAR MATRIX + EXTRA (IRANK-1) WORKING L1 0730
C LOCATIONS. SEE THE COMMENTS IN SUBROUTINE PSIV. L1 0740
C BP AN MM-VECTOR. ITS FIRST IRANK ELEMENTS ARE THE L1 0750
C R.H.S. OF THE TRIANGULAR EQUATIONS AS P*XP=BP. L1 0760
C XP AN MM-VECTOR . ITS FIRST IRANK ELEMENTS ARE THE L1 0770
C SOLUTION OF THE TRIANGULAR EQUATIONS AS P*XP=BP. L1 0780
C UF AN MM-WORKING VECTOR. L1 0790
DOUBLE PRECISION S,SA,SB L1 0800
DIMENSION CT(MM,NN),F(NN),A(MM),GINV(MM,MM),P(MMM) L1 0810
DIMENSION IC(NN),IB(NN),R(NN),T(NN),ALFA(NN),IR(MM) L1 0820
DIMENSION UF(MM),BP(MM),XP(MM),VB(MM) L1 0830
IND=0 L1 0840
TPEPS=2.+EPS L1 0850
NMM=(M*(M+3))/2 L1 0860
IRANK=M L1 0870
ITER=0 L1 0880
DO 1 J=1,N L1 0890
IB(J)=1 L1 0900
1 IC(J)=J L1 0910
DO 3 J=1,M L1 0920
IR(J)=J L1 0930
A(J)=0. L1 0940
DO 2 I=1,M L1 0950
2 GINV(I,J)=0. L1 0960
3 GINV(J,J)=1. L1 0970
IOUT=0 L1 0980
C PART 1 OF THE ALGORITHM. L1 0990
4 IOUT=IOUT+1 L1 1000
IF(IOUT.GT.IRANK) GO TO 16 L1 1010
PIV=0. L1 1020
DO 6 J=IOUT,N L1 1030
ICJ=IC(J) L1 1040
DO 5 I=IOUT,IRANK L1 1050
D=CT(I,ICJ) L1 1060
IF(D.LT.0.0) D=-D L1 1070

```

IF(D.LE.PIV) GO TO 5	L1 1080
LI=I	L1 1090
JIN=ICJ	L1 1100
LJ=J	L1 1110
PIV=D	L1 1120
5 CONTINUE	L1 1130
6 CONTINUE	L1 1140
C DETECTION OF RANK DEFICIENCY.	L1 1150
IF(PIV.GT.EPS) GO TO 7	L1 1160
IRANK=IOUT-1	L1 1170
IND=1	L1 1180
GO TO 16	L1 1190
7 IF(LI.EQ.IOUT) GO TO 10	L1 1200
DO 8 J=1,N	L1 1210
G=CT(LI,J)	L1 1220
CT(LI,J)=CT(IOUT,J)	L1 1230
8 CT(IOUT,J)=G	L1 1240
K=IR(LI)	L1 1250
IR(LI)=IR(IOUT)	L1 1260
IR(IOUT)=K	L1 1270
IF(IOUT.EQ.1) GO TO 10	L1 1280
K=IOUT-1	L1 1290
DO 9 J=1,K	L1 1300
D=GINV(LI,J)	L1 1310
GINV(LI,J)=GINV(IOUT,J)	L1 1320
9 GINV(IOUT,J)=D	L1 1330
C A GAUSS-JORDAN ELIMINATION STEP.	L1 1340
10 PIVOT=CT(IOUT,JIN)	L1 1350
DO 11 J=1,N	L1 1360
11 CT(IOUT,J)=CT(IOUT,J)/PIVOT	L1 1370
DO 12 J=1,IOUT	L1 1380
12 GINV(IOUT,J)=GINV(IOUT,J)/PIVOT	L1 1390
DO 15 I=1,IRANK	L1 1400
IF(I.EQ.IOUT) GO TO 15	L1 1410
D=CT(I,JIN)	L1 1420
DO 13 J=1,N	L1 1430
13 CT(I,J)=CT(I,J)-D*CT(IOUT,J)	L1 1440
DO 14 J=1,IOUT	L1 1450
14 GINV(I,J)=GINV(I,J)-D*GINV(IOUT,J)	L1 1460
15 CONTINUE	L1 1470
ITER=ITER+1	L1 1480
K=IC(LJ)	L1 1490
IC(LJ)=IC(IOUT)	L1 1500
IC(IOUT)=K	L1 1510
GO TO 4	L1 1520
C PART 2 OF THE ALGORITHM.	L1 1530
16 IRANK1=IRANK+1	L1 1540
IRNKM1=IRANK-1	L1 1550
C INITIAL RESIDUALS AND INITIAL BASIC SOLUTION.	L1 1560
DO 17 J=1,IRANK	L1 1570
ICJ=IC(J)	L1 1580
R(ICJ)=0.	L1 1590
17 UF(J)=F(ICJ)	L1 1600
DO 19 J=IRANK1,N	L1 1610
ICJ=IC(J)	L1 1620
S=-F(ICJ)	L1 1630
DO 18 I=1,IRANK	L1 1640
SA=UF(I)	L1 1650
SB=CT(I,ICJ)	L1 1660
18 S=S+SA*SB	L1 1670
R(ICJ)=S	L1 1680
IF(S.GE.0.0) GO TO 19	L1 1690
IB(ICJ)=-1	L1 1700
19 CONTINUE	L1 1710
DO 21 I=1,IRANK	L1 1720
S=0.	L1 1730
DO 20 J=1,N	L1 1740
SA=CT(I,J)	L1 1750
IF(IB(J).EQ.(-1)) SA=-SA	L1 1760
20 S=S+SA	L1 1770
21 VB(I)=S	L1 1780
C INITIALIZING THE TRIANGULAR MATRIX.	L1 1790
DO 22 I=1,NMM	L1 1800
22 P(I)=0.	L1 1810
K=1	L1 1820
KD=IRANK1	L1 1830

DO 23 I=1,IRANK	L1 1840
P(K)=1.	L1 1850
K=K+KD	L1 1860
23 KD=KD-1	L1 1870
C DETERMINE THE VECTOR WHICH LEAVES THE BASIS.	L1 1880
24 IVO=0	L1 1890
CALL PSLV(1,IRANK,MM,MMM,P,VB,XP)	L1 1900
G=1.	L1 1910
DO 27 I=1,IRANK	L1 1920
E=XP(I)	L1 1930
IF(E.LT.(-EPS)) GO TO 25	L1 1940
IF(E.LE.TPEPS) GO TO 27	L1 1950
D=2.0-E	L1 1960
IF(D.GE.G) GO TO 27	L1 1970
IVO=1	L1 1980
GO TO 26	L1 1990
25 D=E	L1 2000
IF(D.GE.G) GO TO 27	L1 2010
IVO=-1	L1 2020
26 G=D	L1 2030
IOUT=I	L1 2040
XB=E	L1 2050
27 CONTINUE	L1 2060
IF(IVO.EQ.0) GO TO 66	L1 2070
C CALCULATION OF ROW (IOUT) IN THE TABLEAU.	L1 2080
IC IOUT=IC(IOUT)	L1 2090
T(IC IOUT)=1.	L1 2100
BXB=XB	L1 2110
DO 28 I=1,IRANK	L1 2120
28 BP(I)=0.	L1 2130
BP(IOUT)=1.	L1 2140
CALL PSLV(2,IRANK,MM,MMM,P,BP,XP)	L1 2150
ALFMX=0.0	L1 2160
DO 31 J=IRANK1,N	L1 2170
ICJ=IC(J)	L1 2180
ALFA(ICJ)=0.	L1 2190
S=0.	L1 2200
DO 29 I=IOUT,IRANK	L1 2210
SA=XP(I)	L1 2220
SB=CT(I,ICJ)	L1 2230
29 S=S+SA*SB	L1 2240
E=S	L1 2250
T(ICJ)=E	L1 2260
IF(E.LT.EPS.AND.E.GT.(-EPS)) GO TO 31	L1 2270
D=R(ICJ)	L1 2280
IF(D.NE.0.0) GO TO 30	L1 2290
D=PREC*PREC*FLOAT(IB(ICJ))	L1 2300
30 ALFA(ICJ)=D/E	L1 2310
GG=ALFA(ICJ)	L1 2320
IF(GG.LT.0.0) GG=-GG	L1 2330
IF(GG.LE.ALFMX) GO TO 31	L1 2340
ALFMX=GG	L1 2350
31 CONTINUE	L1 2360
PIVOTO=1.	L1 2370
ITEST=0	L1 2380
C DETERMINE THE VECTOR WHICH ENTERS THE BASIS.	L1 2390
GG=ALFMX+ALFMX	L1 2400
32 ALFMX=GG	L1 2410
ALFMN=-GG	L1 2420
DO 35 J=IRANK1,N	L1 2430
ICJ=IC(J)	L1 2440
E=ALFA(ICJ)	L1 2450
D=E*FLOAT(IVO)	L1 2460
IF(D.LE.0.0) GO TO 35	L1 2470
IF(IVO.EQ.1) GO TO 33	L1 2480
IF(E.LE.ALFMN) GO TO 35	L1 2490
ALFMN=E	L1 2500
GO TO 34	L1 2510
33 IF(E.GE.ALFMX) GO TO 35	L1 2520
ALFMX=E	L1 2530
34 JIN=J	L1 2540
ITEST=1	L1 2550
35 CONTINUE	L1 2560
IF(ITEST.EQ.1) GO TO 36	L1 2570
C NO FEASIBLE SOLUTION HAS BEEN FOUND.	L1 2580

IND=-1	L1	2590
GO TO 66	L1	2600
36 ICJIN=IC(JIN)	L1	2610
PIVOT=T(ICJIN)	L1	2620
ALPHA=ALFA(ICJIN)	L1	2630
PIVOTN=PIVOT/PIVOTO	L1	2640
IF(XB.GT.TPEPS) GO TO 37	L1	2650
IF(PIVOTN.GT.0.) GO TO 39	L1	2660
GO TO 41	L1	2670
37 DO 38 I=1,IRANK	L1	2680
E=CT(I,ICIOUT)	L1	2690
38 VB(I)=VB(I)-E-E	L1	2700
E=T(ICIOUT)	L1	2710
BXB=BXB-E-E	L1	2720
IB(ICIOUT)=-1	L1	2730
IF(PIVOTN.GT.0.) GO TO 41	L1	2740
39 DO 40 I=1,IRANK	L1	2750
E=CT(I,ICJIN)	L1	2760
40 VB(I)=VB(I)+E+E	L1	2770
E=T(ICJIN)	L1	2780
BXB=BXB+E+E	L1	2790
IB(ICJIN)=1	L1	2800
41 XB=BXB/PIVOT	L1	2810
IF(XB.GE.(-EPS).AND.XB.LE.TPEPS) GO TO 42	L1	2820
ITEST=0	L1	2830
42 ALFA(ICJIN)=0.	L1	2840
IF(ITEST.EQ.1) GO TO 43	L1	2850
PIVOTO=PIVOT	L1	2860
ICIOUT=ICJIN	L1	2870
GO TO 32	L1	2880
43 R(ICJIN)=0.	L1	2890
ITER=ITER+1	L1	2900
IF(IOUT.EQ.IRANK) GO TO 46	L1	2910
C UPDATING MATRIX (P,GINV,VB,CT).	L1	2920
DO 45 J=IOUT,IRNKMI	L1	2930
K=J	L1	2940
K1=K+1	L1	2950
KD=IRANK	L1	2960
L=IC(K)	L1	2970
IC(K)=IC(K1)	L1	2980
IC(K1)=L	L1	2990
DO 44 I=1,K1	L1	3000
P(K)=P(K+1)	L1	3010
K=K+KD	L1	3020
44 KD=KD-1	L1	3030
45 CONTINUE	L1	3040
46 L=IC(IRANK)	L1	3050
IC(IRANK)=IC(JIN)	L1	3060
IC(JIN)=L	L1	3070
K=IRANK	L1	3080
KD=IRANK	L1	3090
DO 47 I=1,IRANK	L1	3100
P(K)=CT(I,ICJIN)	L1	3110
K=K+KD	L1	3120
47 KD=KD-1	L1	3130
IF(IOUT.EQ.IRANK) GO TO 58	L1	3140
DO 57 I=IOUT,IRNKMI	L1	3150
II=I	L1	3160
I1=I+1	L1	3170
K=0	L1	3180
KD=IRANK1	L1	3190
DO 48 J=1,II	L1	3200
K=K+KD	L1	3210
48 KD=KD-1	L1	3220
KK=K	L1	3230
KL=K-KD	L1	3240
L=KL	L1	3250
G=P(K)	L1	3260
D=P(L)	L1	3270
IF(G.LT.0.0) G=-G	L1	3280
IF(D.LT.0.0) D=-D	L1	3290
IF(G.LE.D) GO TO 53	L1	3300
DO 49 J=II,IRANK	L1	3310
E=P(K)	L1	3320
P(K)=P(L)	L1	3330
P(L)=E	L1	3340

```

      K=K+1
49      L=L+1
      J=IR(I)
      IR(I)=IR(I1)
      IR(I1)=J
      DO 50 J=1,N
         E=CT(I,J)
         CT(I,J)=CT(I1,J)
50      CT(I1,J)=E
      DO 51 J=1,IRANK
         E=GINV(I,J)
         GINV(I,J)=GINV(I1,J)
51      GINV(I1,J)=E
      DO 52 J=1,IRANK
         E=GINV(J,I)
         GINV(J,I)=GINV(J,I1)
52      GINV(J,I1)=E
      E=VB(I)
      VB(I)=VB(I1)
      VB(I1)=E
53      E=P(KK)/P(KL)
      IF(E.LT.PREC.AND.E.GT.(-PREC)) GO TO 57
      K=KK
      L=KL
      DO 54 J=II,IRANK
         P(K)=P(K)-E*P(L)
         K=K+1
54      L=L+1
      DO 55 J=1,N
         CT(I1,J)=CT(I1,J)-E*CT(I,J)
55      DO 56 J=1,IRANK
         GINV(I1,J)=GINV(I1,J)-E*GINV(I,J)
56      VB(I1)=VB(I1)-E*VB(I)
57      CONTINUE
58 DO 59 J=1,IRANK
      ICJ=IC(J)
59      UF(J)=F(ICJ)
      IF(ALPHA.GT.(1.0).OR.ALPHA.LT.(-1.0)) GO TO 61
      DO 60 J=IRANK1,N
         ICJ=IC(J)
60      R(ICJ)=R(ICJ)-ALPHA*T(ICJ)
      GO TO 64
61 CALL PSLV(2,IRANK,MM,MMM,P,UF,XP)
      DO 63 J=IRANK1,N
         ICJ=IC(J)
         S=-F(ICJ)
         DO 62 I=1,IRANK
            SA=XP(I)
            SB=CT(I,ICJ)
62      S=S+SA*SB
63      R(ICJ)=S
64 DO 65 J=IRANK1,N
         ICJ=IC(J)
         D=R(ICJ)*FLOAT(IB(ICJ))
         IF(D.GE.0.0) GO TO 65
         R(ICJ)=0.0
65      CONTINUE
      GO TO 24
C CALCULATING THE ANSWER OF THE PROBLEM.
66 CALL PSLV(2,IRANK,MM,MMM,P,UF,VB)
      DO 68 I=1,IRANK
         S=0.
         DO 67 K=1,IRANK
            SA=VB(K)
            SB=GINV(K,I)
67      S=S+SA*SB
         K=IR(I)
68      A(K)=S
         S=0.
         DO 69 J=1,N
            SA=R(J)
            IF(SA.LT.0.0) SA=-SA
69      S=S+SA
         Z=S
         IF(IND.NE.0) RETURN
         E=2.-EPS

```

```

DO 70 I=1,IRANK
  D=XP(I)
  IF(D.LT.EPS.OR.D.GT.E) IND=1
70 CONTINUE
RETURN
END
    
```

```

L1 4110
L1 4120
L1 4130
L1 4140
L1 4150
L1 4160
    
```

```

1.478689 1.113268 0.855048 0.673991 0.541967
0.442372 0.365257 0.304338 0.255453 0.215730
1.027547 0.884824 0.730009 0.598490 0.492769
0.408461 0.340883 0.286251 0.241692 0.205051
0.692608 0.652965 0.577505 0.496487 0.422211
0.357993 0.303718 0.258208 0.220102 0.188148
0.443503 0.452751 0.427474 0.386389 0.341112
0.297441 0.257770 0.222784 0.192393 0.166192
0.251763 0.286357 0.291899 0.279505 0.257866
0.232626 0.207020 0.182721 0.160481 0.140547
0.100448 0.149070 0.173775 0.181324 0.177820
0.167924 0.154817 0.140517 0.126221 0.112594
-0.020740 0.035871 0.072667 0.093931 0.103914
0.106235 0.103674 0.098224 0.091240 0.083607
-0.118431 -0.057358 -0.012925 0.017703 0.037524
0.049293 0.055315 0.057387 0.056853 0.054674
-0.197141 -0.133831 -0.084691 -0.047765 -0.020895
-0.001986 0.010826 0.019096 0.024061 0.026676
-0.260102 -0.196048 -0.144211 -0.103189 -0.071395
-0.047222 -0.029174 -0.015940 -0.006419 0.000282
    
```

```

5. 4. 3. 2. 1. 0.
9. 7. 3. 19. 13. 4.
6. 6. 0. 12. 12. 2.
9. 9. 7. 25. 11. 7.
3. 0. 1. 4. 2. 7.
8. 1. 8. 17. 1. 7.
1. 9. 8. 18. 2. 3.
0. 9. 3. 12. 6. 3.
3. 1. 1. 5. 3. 5.
6. 7. 6. 19. 7. 1.
6. 1. 9. 16. -2. 4.
0. 4. 8. 12. -4. 1.
0. 5. 7. 12. -2. 6.
7. 3. 2. 12. 8. 6.
5. 4. 9. 18. 0. 0.
    
```

```

SUBROUTINE SUNDER
C THIS IS A DUMMY SUBROUTINE. ITS FUNCTION IS TO
C SUPPRESS THE UNDERFLOWS OCCURING IN THE CALCULATION.
RETURN
END
    
```

```

SUN 0010
SUN 0020
SUN 0030
SUN 0040
SUN 0050
    
```

```

SUBROUTINE SETCLK
C THIS IS A DUMMY SUBROUTINE. ITS FUNCTION IS TO START
C READING THE CPU TIME.
RETURN
END
    
```

```

SET 0010
SET 0020
SET 0030
SET 0040
SET 0050
    
```

```

SUBROUTINE RDCLK(TX)
C THIS IS A DUMMY SUBROUTINE. ITS FUNCTION IS TO RECORD
C THE CPU TIME TX IN SECONDS, WHERE TX IS A REAL VARIABLE.
TX=0.00
RETURN
END
    
```

```

RDC 0010
RDC 0020
RDC 0030
RDC 0040
RDC 0050
RDC 0060
    
```

ALGORITHM 552

Solution of the Constrained l_1 Linear Approximation Problem [F4]

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Key Words and Phrases: constrained l_1 approximation, linear programming, simplex method
CR Categories: 5.13, 5.41
Language: Fortran

DESCRIPTION

1. Introduction

Given a $k \times n$ system of linear equations

$$Ax = b, \quad (1)$$

the algorithm calculates an l_1 solution of (1) subject to l linear equality constraints

$$Cx = d \quad (2)$$

and m linear inequality constraints

$$Ex \leq f; \quad (3)$$

i.e., the subroutine determines a column vector x^* which minimizes

$$\|b - Ax\|_1 = \sum_{i=1}^k |b_i - A_i x| \quad (4)$$

subject to the given constraints (2) and (3). (In expression (4), b_i denotes the i th component of b and A_i denotes the i th row of A .)

The method is completely general in the sense that no restrictions are imposed on the ranks of the matrices A , C , and E , or on the signs of the elements of f . Furthermore, if no vector x satisfying (2) and (3) exists, the subroutine detects this and informs the user that the problem is infeasible.

The algorithm can be used to solve the constrained l_1 approximation problem. Suppose that data consisting of k points (t_i, y_i) are to be approximated by a linear approximating function $x_1\phi_1(t) + x_2\phi_2(t) + \dots + x_n\phi_n(t)$, where certain linear constraints are imposed on the parameters x_1, x_2, \dots, x_n . This is equivalent to finding an l_1 solution to the system of equations

$$\sum_{j=1}^n \phi_j(t_i)x_j = y_i, \quad i = 1, 2, \dots, k,$$

subject to the given linear constraints. If the data values y_i contain some large

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errors, then an l_1 approximation may be preferable to an l_2 (least-squares) solution. This aspect of robust regression is discussed, for example, in [5, 7].

The algorithm is a modification of the simplex method applied to the primal formulation of the constrained l_1 problem as a linear program. The modification allows many intermediate simplex vertices to be bypassed, thus reducing the total number of simplex iterations required. A complete description of the method is given in [3]. In the absence of constraints, the algorithm reduces to our earlier unconstrained l_1 algorithm [1], which has proved to be most effective in practice. Indeed, in this case the program which follows produces identical results to the code given in [2].

Following [3], the constrained l_1 problem is posed as the following linear programming problem:

$$\begin{aligned} \text{minimize} \quad & e(u + v) + Me'(u' + v') + Me''v'' \\ \text{subject to} \quad & A(x' - x'') + u - v = b \\ & C(x' - x'') + u' - v' = d \\ & E(x' - x'') + u'' - v'' = f \\ & x', x'' \geq 0, \quad u, v \geq 0, \quad u', v' \geq 0, \quad u'', v'' \geq 0. \end{aligned} \tag{5}$$

In (5), e , e' , and e'' are vectors of 1's of appropriate dimensions and u' , v' , and v'' are artificial vectors with a large positive cost M in the objective function. Our modified simplex method is applied to a condensed tableau of size $(k + l + m + 2) \times (n + 2)$, using the usual two-phase approach in which the artificial part of the objective function is minimized first.

Our algorithm can be used to calculate a weighted constrained l_1 solution, merely by multiplying A_i and b_i by a chosen positive weight before entering the subroutine. Also, as is explained in the comment section of our listing, simple nonnegativity constraints on the variables need not be included explicitly in the constraints (3), since these can be handled most efficiently by assigning the appropriate variables x_j'' in (5) a cost M in the objective function, thus making them artificial variables. Similarly, a nonnegativity constraint on the i th row $b_i - A_i x$ can be handled by making the variable v_i in (5) artificial.

2. Numerical Results

The algorithm has been tested on a variety of problems, several of which are presented in [3]. We include here the following three additional test problems, which have been run on an IBM 370/158 using single-precision arithmetic (approximately 7 decimal digits).

Example 1. Bartels et al. [4] present an example of a spline approximation due to Maurice Cox which is constrained to be convex. This corresponds to an l_1 solution of a 9×7 system of equations whose unknowns are subject to 5 inequality constraints. Our algorithm requires 10 simplex iterations to reproduce (to single-precision accuracy) the results given in [4], where 17 iterations of their penalty function technique are required.

Example 2. This example was devised to demonstrate the ability of our algorithm to handle rank deficient problems. The matrices and vectors are defined as follows:

$$A = \begin{bmatrix} 2 & 0 & 1 & 3 & 1 \\ 7 & 4 & 4 & 15 & 7 \\ 9 & 4 & 7 & 20 & 6 \\ 2 & 2 & 1 & 5 & 3 \\ 9 & 3 & 2 & 14 & 10 \\ 4 & 5 & 0 & 9 & 9 \\ 4 & 4 & 9 & 17 & -1 \\ 1 & 6 & 2 & 9 & 5 \end{bmatrix}, \quad b = \begin{bmatrix} 7 \\ 4 \\ 7 \\ 4 \\ 0 \\ 4 \\ 9 \\ 6 \end{bmatrix},$$

$$C = \begin{bmatrix} 0 & 4 & 5 & 9 & -1 \\ 3 & 2 & 7 & 12 & -2 \\ 3 & 6 & 12 & 21 & -3 \end{bmatrix}, \quad d = \begin{bmatrix} 5 \\ 1 \\ 6 \end{bmatrix},$$

$$E = \begin{bmatrix} 0 & 3 & 6 & 9 & -3 \\ 6 & 2 & 4 & 12 & 4 \end{bmatrix}, \quad f = \begin{bmatrix} 5 \\ 6 \end{bmatrix}.$$

The 8×5 matrix A is of rank 3. The 3×5 matrix C is of rank 2, which causes one artificial vector to remain in the basis at the end of phase I. Our algorithm includes an adaptation of a standard device (e.g., see Hadley [6, p. 153]) to handle such situations. The solution is obtained with no difficulty in 3 iterations with $\text{ERROR} = 26.148$.

Example 3. Compute a best l_1 approximation by $p_5(t) = \sum_{j=1}^6 x_j t^{j-1}$ to the function $y(t) = \min(e^t, e^{1/2})$ on the 101 points defined by $t = 0(0.01)1$, subject to the interpolation condition that $p_5(t) = y(t)$ for $t = 0, 0.5$, and 1. This gives rise to a 98×6 overdetermined system with 3 equality constraints. The algorithm requires 8 iterations to solve the problem (with $\text{ERROR} = 2.8513$), and it requires 13 iterations to solve the 101×6 overdetermined system *without* the 3 constraints (with $\text{ERROR} = 1.1473$).

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ALGORITHM

C SAMPLE DRIVER PROGRAM FOR SUBROUTINE CL1.	
C	MAN 20
C THIS PROGRAM SOLVES A K BY N OVERDETERMINED SYSTEM	MAN 30
C	MAN 40
C AX=B	MAN 50
C	MAN 60
C IN THE L1 SENSE SUBJECT TO L EQUALITY CONSTRAINTS	MAN 70
C	MAN 80
C CX=D	MAN 90
C	MAN 100
C AND M INEQUALITY CONSTRAINTS	MAN 110
C	MAN 120
C EX.LE.F.	MAN 130
C	MAN 140
C COMPLETE DETAILS OF THE PARAMETERS MAY BE	MAN 150
C FOUND IN THE DOCUMENTATION OF THE SUBROUTINE.	MAN 160
C	MAN 170
C THE ARRAYS ARE CURRENTLY DIMENSIONED TO ALLOW PROBLEMS	MAN 180
C FOR WHICH K+L+M .LE. 100, N .LE. 10.	MAN 190
C	MAN 200
C THE PROGRAM MAY BE TESTED ON THE FOLLOWING DATA.	MAN 210
C	MAN 220
C K = 8	MAN 230
C L = 3	MAN 240
C M = 2	MAN 250
C N = 5	MAN 260
C	MAN 270
C Q = 2 0 1 3 1 7	MAN 280
C 7 4 4 15 7 4	MAN 290
C 9 4 7 20 6 7	MAN 300

C	2 2 1 5 3 4	MAN	310
C	9 3 2 14 10 0	MAN	320
C	4 5 0 9 9 4	MAN	330
C	4 4 9 17 -1 9	MAN	340
C	1 6 2 9 5 6	MAN	350
C	0 4 5 9 -1 5	MAN	360
C	3 2 7 12 -2 1	MAN	370
C	3 6 12 21 -3 6	MAN	380
C	0 3 6 9 -3 5	MAN	390
C	6 2 4 12 4 6	MAN	400
C		MAN	410
C	KODE = 0	MAN	420
C	TOLER = 1.E-5	MAN	430
C	ITER = 130	MAN	440
C		MAN	450
C		MAN	460
C	DIMENSION Q(102,12), X(12), RES(100), CU(2,110)	MAN	470
C	INTEGER IU(2,110), S(100)	MAN	480
C	DATA KLMD, KLM2D, NKLMD, N2D /100,102,110,12/	MAN	490
C	INPUT DATA.	MAN	500
C	READ (5,99999) K, L, M, N, KODE, TOLER, ITER	MAN	510
C	KLM = K + L + M	MAN	520
C	N1 = N + 1	MAN	530
C	DO 10 I=1,KLM	MAN	540
C	READ (5,99998) (Q(I,J),J=1,N1)	MAN	550
C	WRITE (6,99994) (Q(I,J),J=1,N1)	MAN	560
C	10 CONTINUE	MAN	570
C	CALL CL1(K, L, M, N, KLMD, KLM2D, NKLMD, N2D, Q,	MAN	580
C	* KODE, TOLER, ITER, X, RES, ERROR, CU, IU, S)	MAN	590
C	OUTPUT KODE, ITERATION COUNT AND ERROR NORM.	MAN	600
C	WRITE (6,99997) KODE, ITER, ERROR	MAN	610
C	OUTPUT SOLUTION VECTOR.	MAN	620
C	WRITE (6,99996) (I,X(I),I=1,N)	MAN	630
C	OUTPUT RESIDUAL ERROR AT EACH POINT.	MAN	640
C	WRITE (6,99995) (I,RES(I),I=1,KLM)	MAN	650
C	STOP	MAN	660
C	99999 FORMAT (5I3, E10.0, I3)	MAN	670
C	99998 FORMAT (8F3.0)	MAN	680
C	99997 FORMAT (16H KODE,ITER,ERROR, 2I10, E18.7)	MAN	690
C	99996 FORMAT (4H SOL, I5, E18.7)	MAN	700
C	99995 FORMAT (6H ERROR, I5, E18.7)	MAN	710
C	99994 FORMAT (2H , 8F5.0)	MAN	720
C	END	MAN	730
C	SUBROUTINE CL1(K, L, M, N, KLMD, KLM2D, NKLMD, N2D,	CL1	10
C	* Q, KODE, TOLER, ITER, X, RES, ERROR, CU, IU, S)	CL1	20
C	THIS SUBROUTINE USES A MODIFICATION OF THE SIMPLEX	CL1	30
C	METHOD OF LINEAR PROGRAMMING TO CALCULATE AN LI SOLUTION	CL1	40
C	TO A K BY N SYSTEM OF LINEAR EQUATIONS	CL1	50
C	AX=B	CL1	60
C	SUBJECT TO L LINEAR EQUALITY CONSTRAINTS	CL1	70
C	CX=D	CL1	80
C	AND M LINEAR INEQUALITY CONSTRAINTS	CL1	90
C	EX.LE.F.	CL1	100
C	DESCRIPTION OF PARAMETERS	CL1	110
C	K NUMBER OF ROWS OF THE MATRIX A (K.GE.1).	CL1	120
C	L NUMBER OF ROWS OF THE MATRIX C (L.GE.0).	CL1	130
C	M NUMBER OF ROWS OF THE MATRIX E (M.GE.0).	CL1	140
C	N NUMBER OF COLUMNS OF THE MATRICES A,C,E (N.GE.1).	CL1	150
C	KLMD SET TO AT LEAST K+L+M FOR ADJUSTABLE DIMENSIONS.	CL1	160
C	KLM2D SET TO AT LEAST K+L+M+2 FOR ADJUSTABLE DIMENSIONS.	CL1	170
C	NKLMD SET TO AT LEAST N+K+L+M FOR ADJUSTABLE DIMENSIONS.	CL1	180
C	N2D SET TO AT LEAST N+2 FOR ADJUSTABLE DIMENSIONS	CL1	190
C	Q TWO DIMENSIONAL REAL ARRAY WITH KLM2D ROWS AND	CL1	200
C	AT LEAST N2D COLUMNS.	CL1	210
C	ON ENTRY THE MATRICES A,C AND E, AND THE VECTORS	CL1	220
C	B,D AND F MUST BE STORED IN THE FIRST K+L+M ROWS	CL1	230
C	AND N+1 COLUMNS OF Q AS FOLLOWS	CL1	240
C	A B	CL1	250
C	Q = C D	CL1	260
C	E F	CL1	270
C	THESE VALUES ARE DESTROYED BY THE SUBROUTINE.	CL1	280
C	KODE A CODE USED ON ENTRY TO, AND EXIT	CL1	290
C	FROM, THE SUBROUTINE.	CL1	300
C	ON ENTRY, THIS SHOULD NORMALLY BE SET TO 0.	CL1	310


```

C      HOWEVER, IF CERTAIN NONNEGATIVITY CONSTRAINTS          CL1 320
C      ARE TO BE INCLUDED IMPLICITLY, RATHER THAN             CL1 330
C      EXPLICITLY IN THE CONSTRAINTS EX.LE.F, THEN KODE      CL1 340
C      SHOULD BE SET TO 1, AND THE NONNEGATIVITY             CL1 350
C      CONSTRAINTS INCLUDED IN THE ARRAYS X AND                CL1 360
C      RES (SEE BELOW).                                       CL1 370
C      ON EXIT, KODE HAS ONE OF THE                            CL1 380
C      FOLLOWING VALUES                                       CL1 390
C      0- OPTIMAL SOLUTION FOUND,                              CL1 400
C      1- NO FEASIBLE SOLUTION TO THE                          CL1 410
C      CONSTRAINTS,                                           CL1 420
C      2- CALCULATIONS TERMINATED                             CL1 430
C      PREMATURELY DUE TO ROUNDING ERRORS,                    CL1 440
C      3- MAXIMUM NUMBER OF ITERATIONS REACHED.              CL1 450
C TOLER  A SMALL POSITIVE TOLERANCE. EMPIRICAL               CL1 460
C      EVIDENCE SUGGESTS TOLER = 10**(-D*2/3),                CL1 470
C      WHERE D REPRESENTS THE NUMBER OF DECIMAL                CL1 480
C      DIGITS OF ACCURACY AVAILABLE. ESSENTIALLY,             CL1 490
C      THE SUBROUTINE CANNOT DISTINGUISH BETWEEN ZERO         CL1 500
C      AND ANY QUANTITY WHOSE MAGNITUDE DOES NOT EXCEED      CL1 510
C      TOLER. IN PARTICULAR, IT WILL NOT PIVOT ON ANY        CL1 520
C      NUMBER WHOSE MAGNITUDE DOES NOT EXCEED TOLER.         CL1 530
C ITER   ON ENTRY ITER MUST CONTAIN AN UPPER BOUND ON        CL1 540
C      THE MAXIMUM NUMBER OF ITERATIONS ALLOWED.              CL1 550
C      A SUGGESTED VALUE IS 10*(K+L+M). ON EXIT ITER         CL1 560
C      GIVES THE NUMBER OF SIMPLEX ITERATIONS.                CL1 570
C X      ONE DIMENSIONAL REAL ARRAY OF SIZE AT LEAST N2D.     CL1 580
C      ON EXIT THIS ARRAY CONTAINS A                           CL1 590
C      SOLUTION TO THE L1 PROBLEM. IF KODE=1                   CL1 600
C      ON ENTRY, THIS ARRAY IS ALSO USED TO INCLUDE           CL1 610
C      SIMPLE NONNEGATIVITY CONSTRAINTS ON THE                 CL1 620
C      VARIABLES. THE VALUES -1, 0, OR 1                     CL1 630
C      FOR X(J) INDICATE THAT THE J-TH VARIABLE                CL1 640
C      IS RESTRICTED TO BE .LE.0, UNRESTRICTED,              CL1 650
C      OR .GE.0 RESPECTIVELY.                                  CL1 660
C RES    ONE DIMENSIONAL REAL ARRAY OF SIZE AT LEAST KLMD.    CL1 670
C      ON EXIT THIS CONTAINS THE RESIDUALS B-AX                CL1 680
C      IN THE FIRST K COMPONENTS, D-CX IN THE                  CL1 690
C      NEXT L COMPONENTS (THESE WILL BE =0), AND                CL1 700
C      F-EX IN THE NEXT M COMPONENTS. IF KODE=1 ON            CL1 710
C      ENTRY, THIS ARRAY IS ALSO USED TO INCLUDE SIMPLE        CL1 720
C      NONNEGATIVITY CONSTRAINTS ON THE RESIDUALS             CL1 730
C      B-AX. THE VALUES -1, 0, OR 1 FOR RES(I)                CL1 740
C      INDICATE THAT THE I-TH RESIDUAL (1.LE.I.LE.K) IS      CL1 750
C      RESTRICTED TO BE .LE.0, UNRESTRICTED, OR .GE.0        CL1 760
C      RESPECTIVELY.                                          CL1 770
C ERROR  ON EXIT, THIS GIVES THE MINIMUM SUM OF               CL1 780
C      ABSOLUTE VALUES OF THE RESIDUALS.                      CL1 790
C CU     A TWO DIMENSIONAL REAL ARRAY WITH TWO ROWS AND      CL1 800
C      AT LEAST NKLMD COLUMNS USED FOR WORKSPACE.           CL1 810
C IU     A TWO DIMENSIONAL INTEGER ARRAY WITH TWO ROWS AND    CL1 820
C      AT LEAST NKLMD COLUMNS USED FOR WORKSPACE.           CL1 830
C S      INTEGER ARRAY OF SIZE AT LEAST KLMD, USED FOR        CL1 840
C      WORKSPACE.                                              CL1 850
C      CL1 860
C IF YOUR FORTRAN COMPILER PERMITS A SINGLE COLUMN OF A TWO  CL1 870
C DIMENSIONAL ARRAY TO BE PASSED TO A ONE DIMENSIONAL ARRAY  CL1 880
C THROUGH A SUBROUTINE CALL, CONSIDERABLE SAVINGS IN         CL1 890
C EXECUTION TIME MAY BE ACHIEVED THROUGH THE USE OF THE     CL1 900
C FOLLOWING SUBROUTINE, WHICH OPERATES ON COLUMN VECTORS.    CL1 910
C      SUBROUTINE COL(V1, V2, XMLT, NOTROW, K)                  CL1 920
C THIS SUBROUTINE ADDS TO THE VECTOR V1 A MULTIPLE OF THE    CL1 930
C VECTOR V2 (ELEMENTS 1 THROUGH K EXCLUDING NOTROW).         CL1 940
C      DIMENSION V1(K), V2(K)                                  CL1 950
C      KEND = NOTROW - 1                                       CL1 960
C      KSTART = NOTROW + 1                                      CL1 970
C      IF (KEND .LT. 1) GO TO 20                                CL1 980
C      DO 10 I=1,KEND                                           CL1 990
C          V1(I) = V1(I) + XMLT*V2(I)                            CL1 1000
C 10 CONTINUE                                                  CL1 1010
C      IF(KSTART .GT. K) GO TO 40                                CL1 1020
C 20 DO 30 I=KSTART,K                                           CL1 1030
C          V1(I) = V1(I) + XMLT*V2(I)                            CL1 1040
C 30 CONTINUE                                                  CL1 1050
C 40 RETURN                                                    CL1 1060
C      END                                                       CL1 1070
C SEE COMMENTS FOLLOWING STATEMENT LABELLED 440 FOR

```

```

C INSTRUCTIONS ON THE IMPLEMENTATION OF THIS MODIFICATION.
  DOUBLE PRECISION SUM
  DOUBLE PRECISION DBLE
  REAL Q, X, Z, CU, SN, ZU, ZV, CUV, RES, XMAX, XMIN,
  * ERROR, PIVOT, TOLER, TPivot
  REAL ABS
  INTEGER I, J, K, L, M, N, S, IA, II, IN, IU, JS, KK,
  * NK, N1, N2, JMN, JPN, KLM, NKL, NK1, N2D, IIMN,
  * IOU, ITER, KLMD, KLM1, KLM2, KODE, NKLM, NKL1,
  * KLM2D, MAXIT, NKLMD, IPHASE, KFORCE, IINEG
  INTEGER IABS
  DIMENSION Q(KLM2D,N2D), X(N2D), RES(KLMD),
  * CU(2,NKLMD), IU(2,NKLMD), S(KLMD)
C
C INITIALIZATION.
C
  MAXIT = ITER
  N1 = N + 1
  N2 = N + 2
  NK = N + K
  NK1 = NK + 1
  NKL = NK + L
  NKL1 = NKL + 1
  KLM = K + L + M
  KLM1 = KLM + 1
  KLM2 = KLM + 2
  NKLM = N + KLM
  KFORCE = 1
  ITER = 0
  JS = 1
  IA = 0
C SET UP LABELS IN Q.
DO 10 J=1,N
  Q(KLM2,J) = J
10 CONTINUE
DO 30 I=1,KLM
  Q(I,N2) = N + I
  IF (Q(I,N1).GE.0.) GO TO 30
  DO 20 J=1,N2
    Q(I,J) = -Q(I,J)
20 CONTINUE
30 CONTINUE
C SET UP PHASE 1 COSTS.
IPHASE = 2
DO 40 J=1,NKLM
  CU(1,J) = 0.
  CU(2,J) = 0.
  IU(1,J) = 0
  IU(2,J) = 0
40 CONTINUE
IF (L.EQ.0) GO TO 60
DO 50 J=NK1,NKL
  CU(1,J) = 1.
  CU(2,J) = 1.
  IU(1,J) = 1
  IU(2,J) = 1
50 CONTINUE
IPHASE = 1
60 IF (M.EQ.0) GO TO 80
DO 70 J=NKL1,NKLM
  CU(2,J) = 1.
  IU(2,J) = 1
  JMN = J - N
  IF (Q(JMN,N2).LT.0.) IPHASE = 1
70 CONTINUE
80 IF (KODE.EQ.0) GO TO 150
DO 110 J=1,N
  IF (X(J)) 90, 110, 100
90 CU(1,J) = 1.
  IU(1,J) = 1
  GO TO 110
100 CU(2,J) = 1.
  IU(2,J) = 1
110 CONTINUE
DO 140 J=1,K
  JPN = J + N

```

```

CL1 1080
CL1 1090
CL1 1100
CL1 1110
CL1 1120
CL1 1130
CL1 1140
CL1 1150
CL1 1160
CL1 1170
CL1 1180
CL1 1190
CL1 1200
CL1 1210
CL1 1220
CL1 1230
CL1 1240
CL1 1250
CL1 1260
CL1 1270
CL1 1280
CL1 1290
CL1 1300
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CL1 1330
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CL1 1360
CL1 1370
CL1 1380
CL1 1390
CL1 1400
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CL1 1500
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CL1 1600
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CL1 1760
CL1 1770
CL1 1780
CL1 1790
CL1 1800
CL1 1810
CL1 1820
CL1 1830

```

```

120     IF (RES(J)) 120, 140, 130
120     CU(1,JPN) = 1.
120     IU(1,JPN) = 1
120     IF (Q(J,N2).GT.0.0) IPHASE = 1
120     GO TO 140
130     CU(2,JPN) = 1.
130     IU(2,JPN) = 1
130     IF (Q(J,N2).LT.0.0) IPHASE = 1
140     CONTINUE
150     IF (IPHASE.EQ.2) GO TO 500
C COMPUTE THE MARGINAL COSTS.
160     DO 200 J=JS,N1
160     SUM = 0.D0
160     DO 190 I=1,KLM
160     II = Q(I,N2)
160     IF (II.LT.0) GO TO 170
160     Z = CU(1,II)
160     GO TO 180
170     IINEG = -II
170     Z = CU(2,IINEG)
180     SUM = SUM + DBLE(Q(I,J))*DBLE(Z)
190     CONTINUE
190     Q(KLML,J) = SUM
200     CONTINUE
200     DO 230 J=JS,N
200     II = Q(KLM2,J)
200     IF (II.LT.0) GO TO 210
200     Z = CU(1,II)
200     GO TO 220
210     IINEG = -II
210     Z = CU(2,IINEG)
220     Q(KLML,J) = Q(KLML,J) - Z
230     CONTINUE
C DETERMINE THE VECTOR TO ENTER THE BASIS.
240     XMAX = 0.
240     IF (JS.GT.N) GO TO 490
240     DO 280 J=JS,N
240     ZU = Q(KLML,J)
240     II = Q(KLM2,J)
240     IF (II.GT.0) GO TO 250
240     II = -II
240     ZV = ZU
240     ZU = -ZU - CU(1,II) - CU(2,II)
240     GO TO 260
250     ZV = -ZU - CU(1,II) - CU(2,II)
260     IF (KFORCE.EQ.1 .AND. II.GT.N) GO TO 280
260     IF (IU(1,II).EQ.1) GO TO 270
260     IF (ZU.LE.XMAX) GO TO 270
260     XMAX = ZU
260     IN = J
270     IF (IU(2,II).EQ.1) GO TO 280
270     IF (ZV.LE.XMAX) GO TO 280
270     XMAX = ZV
270     IN = J
280     CONTINUE
280     IF (XMAX.LE.TOLER) GO TO 490
280     IF (Q(KLML,IN).EQ.XMAX) GO TO 300
280     DO 290 I=1,KLM2
280     Q(I,IN) = -Q(I,IN)
290     CONTINUE
290     Q(KLML,IN) = XMAX
C DETERMINE THE VECTOR TO LEAVE THE BASIS.
300     IF (IPHASE.EQ.1 .OR. IA.EQ.0) GO TO 330
300     XMAX = 0.
300     DO 310 I=1,IA
300     Z = ABS(Q(I,IN))
300     IF (Z.LE.XMAX) GO TO 310
300     XMAX = Z
300     IOUT = I
310     CONTINUE
310     IF (XMAX.LE.TOLER) GO TO 330
310     DO 320 J=1,N2
310     Z = Q(IA,J)
310     Q(IA,J) = Q(IOUT,J)
310     Q(IOUT,J) = Z
320     CONTINUE

```

```

CL1 1840
CL1 1850
CL1 1860
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CL1 1880
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CL1 1930
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CL1 1950
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CL1 1970
CL1 1980
CL1 1990
CL1 2000
CL1 2010
CL1 2020
CL1 2030
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CL1 2060
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CL1 2100
CL1 2110
CL1 2120
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CL1 2170
CL1 2180
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CL1 2570
CL1 2580
CL1 2590

```

```

IOUT = IA
IA = IA - 1
PIVOT = Q(IOUT,IN)
GO TO 420
330 KK = 0
DO 340 I=1,KLM
  Z = Q(I,IN)
  IF (Z.LE.TOLER) GO TO 340
  KK = KK + 1
  RES(KK) = Q(I,N1)/Z
  S(KK) = I
340 CONTINUE
350 IF (KK.GT.0) GO TO 360
  KODE = 2
  GO TO 590
360 XMIN = RES(1)
  IOUT = S(1)
  J = 1
  IF (KK.EQ.1) GO TO 380
  DO 370 I=2,KK
    IF (RES(I).GE.XMIN) GO TO 370
    J = I
    XMIN = RES(I)
    IOUT = S(I)
370 CONTINUE
  RES(J) = RES(KK)
  S(J) = S(KK)
380 KK = KK - 1
  PIVOT = Q(IOUT,IN)
  II = Q(IOUT,N2)
  IF (IPHASE.EQ.1) GO TO 400
  IF (II.LT.0) GO TO 390
  IF (IU(2,II).EQ.1) GO TO 420
  GO TO 400
390 IINEG = -II
  IF (IU(1,IINEG).EQ.1) GO TO 420
400 II = IABS(II)
  CUV = CU(1,II) + CU(2,II)
  IF (Q(KLM1,IN)-PIVOT*CUV.LE.TOLER) GO TO 420
C BYPASS INTERMEDIATE VERTICES.
  DO 410 J=JS,N1
    Z = Q(IOUT,J)
    Q(KLM1,J) = Q(KLM1,J) - Z*CUV
    Q(IOUT,J) = -Z
410 CONTINUE
  Q(IOUT,N2) = -Q(IOUT,N2)
  GO TO 350
C GAUSS-JORDAN ELIMINATION.
420 IF (ITER.LT.MAXIT) GO TO 430
  KODE = 3
  GO TO 590
430 ITER = ITER + 1
  DO 440 J=JS,N1
    IF (J.NE.IN) Q(IOUT,J) = Q(IOUT,J)/PIVOT
440 CONTINUE
C IF PERMITTED, USE SUBROUTINE COL OF THE DESCRIPTION
C SECTION AND REPLACE THE FOLLOWING SEVEN STATEMENTS DOWN
C TO AND INCLUDING STATEMENT NUMBER 460 BY..
C DO 460 J=JS,N1
C IF(J.EQ.IN) GO TO 460
C Z = -Q(IOUT,J)
C CALL COL(Q(1,J), Q(1,IN), Z, IOUT, KLM1)
C 460 CONTINUE
  DO 460 J=JS,N1
    IF (J.EQ.IN) GO TO 460
    Z = -Q(IOUT,J)
    DO 450 I=1,KLM1
      IF (I.NE.IOUT) Q(I,J) = Q(I,J) + Z*Q(I,IN)
450 CONTINUE
460 CONTINUE
  TPIVOT = -PIVOT
  DO 470 I=1,KLM1
    IF (I.NE.IOUT) Q(I,IN) = Q(I,IN)/TPIVOT
470 CONTINUE
  Q(IOUT,IN) = 1./PIVOT
  Z = Q(IOUT,N2)

```

```

CL1 2600
CL1 2610
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CL1 2690
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CL1 2990
CL1 3000
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CL1 3020
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CL1 3100
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CL1 3270
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CL1 3290
CL1 3300
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CL1 3330
CL1 3340
CL1 3350

```

Q(IOUT,N2) = Q(KLM2,IN)	CL1 3360
Q(KLM2,IN) = Z	CL1 3370
II = ABS(Z)	CL1 3380
IF (IU(1,II).EQ.0 .OR. IU(2,II).EQ.0) GO TO 240	CL1 3390
DO 480 I=1,KLM2	CL1 3400
Z = Q(I,IN)	CL1 3410
Q(I,IN) = Q(I,JS)	CL1 3420
Q(I,JS) = Z	CL1 3430
480 CONTINUE	CL1 3440
JS = JS + 1	CL1 3450
GO TO 240	CL1 3460
C TEST FOR OPTIMALITY.	CL1 3470
490 IF (KFORCE.EQ.0) GO TO 580	CL1 3480
IF (IPHASE.EQ.1 .AND. Q(KLM1,N1).LE.TOLER) GO TO 500	CL1 3490
KFORCE = 0	CL1 3500
GO TO 240	CL1 3510
C SET UP PHASE 2 COSTS.	CL1 3520
500 IPHASE = 2	CL1 3530
DO 510 J=1,NKLM	CL1 3540
CU(1,J) = 0.	CL1 3550
CU(2,J) = 0.	CL1 3560
510 CONTINUE	CL1 3570
DO 520 J=N1,NK	CL1 3580
CU(1,J) = 1.	CL1 3590
CU(2,J) = 1.	CL1 3600
520 CONTINUE	CL1 3610
DO 560 I=1,KLM	CL1 3620
II = Q(I,N2)	CL1 3630
IF (II.GT.0) GO TO 530	CL1 3640
II = -II	CL1 3650
IF (IU(2,II).EQ.0) GO TO 560	CL1 3660
CU(2,II) = 0.	CL1 3670
GO TO 540	CL1 3680
530 IF (IU(1,II).EQ.0) GO TO 560	CL1 3690
CU(1,II) = 0.	CL1 3700
540 IA = IA + 1	CL1 3710
DO 550 J=1,N2	CL1 3720
Z = Q(IA,J)	CL1 3730
Q(IA,J) = Q(I,J)	CL1 3740
Q(I,J) = Z	CL1 3750
550 CONTINUE	CL1 3760
560 CONTINUE	CL1 3770
GO TO 160	CL1 3780
570 IF (Q(KLM1,N1).LE.TOLER) GO TO 500	CL1 3790
KODE = 1	CL1 3800
GO TO 590	CL1 3810
580 IF (IPHASE.EQ.1) GO TO 570	CL1 3820
C PREPARE OUTPUT.	CL1 3830
KODE = 0	CL1 3840
590 SUM = 0.D0	CL1 3850
DO 600 J=1,N	CL1 3860
X(J) = 0.	CL1 3870
600 CONTINUE	CL1 3880
DO 610 I=1,KLM	CL1 3890
RES(I) = 0.	CL1 3900
610 CONTINUE	CL1 3910
DO 640 I=1,KLM	CL1 3920
II = Q(I,N2)	CL1 3930
SN = 1.	CL1 3940
IF (II.GT.0) GO TO 620	CL1 3950
II = -II	CL1 3960
SN = -1.	CL1 3970
620 IF (II.GT.N) GO TO 630	CL1 3980
X(II) = SN*Q(I,N1)	CL1 3990
GO TO 640	CL1 4000
630 IIMN = II - N	CL1 4010
RES(IIMN) = SN*Q(I,N1)	CL1 4020
IF (II.GE.N1 .AND. II.LE.NK) SUM = SUM +	CL1 4030
* DBLE(Q(I,N1))	CL1 4040
640 CONTINUE	CL1 4050
ERROR = SUM	CL1 4060
RETURN	CL1 4070
END	CL1 4080

8 3 2 5 0 1.E-5130
2 0 1 3 1 7
7 4 4 15 7 4
9 4 7 20 6 7
2 2 1 5 3 4
9 3 2 14 10 0
4 5 0 9 9 4
4 4 9 17 -1 9
1 6 2 9 5 6
0 4 5 9 -1 5
3 2 7 12 -2 1
3 6 12 21 -3 6
0 3 6 9 -3 5
6 2 4 12 4 6

ALGORITHM 553

M3RK, An Explicit Time Integrator for Semidiscrete Parabolic Equations [D3]

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Mathematical Center, Amsterdam, The Netherlands

Key Words and Phrases: parabolic partial differential equations, semidiscretization, explicit time integrator

CR Categories: 5.17

Language: Fortran

DESCRIPTION

This algorithm is a complement to [2] where it is explained and described.

REFERENCES

1. RYDER, B.G. The PFORT verifier. *Softw. Pract. Exper.* 4 (1974), 359-378.
2. VERWER, J.G. An implementation of a class of stabilized explicit methods for the time integration of parabolic equations. *ACM Trans. Math. Softw.* 6, 2 (June 1980), 188-205.

ALGORITHM

```

C      PROGRAM TEST1(OUTPUT,TAPE6=OUTPUT)
C
C TEST PROGRAM 1 FOR J.G.VERWER'S ALGORITHM M3RK, AN EXPLICIT TIME-
C INTEGRATOR FOR SEMI-DISCRETE PARABOLIC EQUATIONS.
C THE TEST PROBLEM IS A SYSTEM OF 2 ONE-DIMENSIONAL PARABOLIC EQUA-
C TIONS WHICH IS DISCUSSED IN SECTION 5.1 OF J.G.VERWER'S COMPANION
C PAPER "AN IMPLEMENTATION OF A CLASS OF STABILIZED, EXPLICIT ME-
C THODS FOR THE TIME INTEGRATION OF PARABOLIC EQUATIONS".
C THE TEST CONSISTS OF THE INTEGRATION OF THE SYSTEM OF 122 ODE'S
C WHICH IS CALLED SYSTEM II IN SECTION 5.1 OF THE COMPANION PAPER
C (EQUATIONS (5.3) WITH 61 GRID POINTS FOR THE GALERKIN DISCRETI-
C ZATION). IN THE PRESENT TEST THE TOLERANCE PARAMETER TOL=1.0E-4.
C
C *****
C
C ARRAY DECLARATIONS FOR M3RK. THE ARRAY U IS THE INPUT ARRAY FOR
C THE INITIAL VECTOR AND UOUT IS THE OUTPUT ARRAY FOR THE COMPUTED
C SOLUTION VECTOR.
C DER IS THE SUBROUTINE DEFINING THE SYSTEM OF ODE'S.
C THE ARRAY X WILL CONTAIN THE GRID POINTS FOR THE GALERKIN DISCRE-
C TIZATION AND IS USED BY DER.
C
      DIMENSION INFO(15),U(122),U1(122),U2(122),UOUT(122),DU(122)
      1,DU1(122),SIGMA(2)
      EXTERNAL DER
      COMMON /GRID/ X(61)
C
C DEFINITION AND PRINTING OF THE GRID POINTS X(I).THE NUMBER OF

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```

C GRID POINTS CAN BE CHANGED BY REDEFINING NPPTS AND BY ADJUSTING
C THE LENGTH OF THE ARRAYS.
C
  NPPTS=61
  DO 10 I=1, NPPTS
10 X(I)=(FLOAT(I)-1.)/(FLOAT(NPPTS)-1.)
  WRITE(6, 12)
12 FORMAT(7H1 GRID=)
  DO 13 I=1, NPPTS
13 WRITE(6, 14)X(I)
14 FORMAT(F12.5)
C
C DEFINITION AND PRINTING OF THE INITIAL VECTOR FOR THE ODE'S.
C
  WRITE(6, 20)
20 FORMAT(17H1 INITIAL VALUES=)
  WRITE(6, 55)
  DO 30 I = 1, NPPTS
  U(I) = 1.0
  IN=I+NPPTS
  U(IN) = .0
  WRITE(6, 70)X(I), U(I), U(IN)
30 CONTINUE
C
C DEFINITION OF PARAMETERS FOR M3RK. NEQN GIVES THE NUMBER OF COM-
C PONENTS OF THE SYSTEM OF ODE'S. T IS THE INDEPENDENT VARIABLE.
C THE MEANING OF THE OTHER PARAMETERS IS CLEARLY EXPLAINED IN THE
C PROLOGUE OF COMMENTS OF M3RK.
C
  NEQN=NPPTS*2
  T=.0
  TOL=1.E-4
  INFO(1)=0
  INFO(2)=2
  INFO(3)=5000
C
C INITIAL AND SUBSEQUENT CALLS OF M3RK(SEE SECTION 5.1 OF THE
C COMPANION PAPER). TE DEFINES THE POINT WHERE OUTPUT IS DESIRED.
C
  DO 80 I=1, 6
  TE=.01
  IF(1.EQ.2)TE=.1
  IF(1.EQ.3)TE=1.
  IF(1.EQ.4)TE=5.
  IF(1.EQ.5)TE=10.
  IF(1.EQ.6)TE=20.
  CALL M3RK(T, TE, NEQN, H, HMIN, SIGMA, TOL,
+         DER, U, U1, U2, UOUT, DU, DU1, IFLAG, INFO)
C
C PRINTING OF THE OUTPUT POINT, OF THE ERROR FLAG, OF THE ESTIMATED
C SPECTRAL RADIUS, OF THE ARRAY INFO, AND OF THE COMPUTED SOLUTION.
C
  WRITE(6, 40)TE, IFLAG, SIGMA(1)
40 FORMAT(14H1 ENDPPOINT =, E10.4/10H IFLAG= , I2/8H SIGMA=
  1, E13.7)
  WRITE(6, 50)(INFO(KK), KK=1, 15)
50 FORMAT(8H INFO= , 15I6/)
  WRITE(6, 55)
55 FORMAT(5X, 1HX, 17X, 1HU, 19X, 1HV/)
  DO 60 J=1, NPPTS
  JN=J+NPPTS
60 WRITE(6, 70)X(J), UOUT(J), UOUT(JN)
70 FORMAT(F9.5, 2F20.7)
80 CONTINUE
  STOP
  END

  SUBROUTINE DER(N, Y)
C
C SUBROUTINE DER DEFINES THE SYSTEM OF ODE'S BEING INTEGRATED.
C THE ARRAY U IS A WORK ARRAY.
C
  DIMENSION Y(N), U(122)
  COMMON /GRID/ X(61)
  FI(Z)=EXP(Z*5.73)-EXP(-11.46*Z)

```



```

      DO 10 I=1,N
10  U(I)=Y(I)
      EPS=.143
      P=.1743
      Y(1)=-2.*EPS*P*(7.*U(1)-8.*U(2)+U(3))/(X(3)-X(1))**2-FI(U(1)-U(62)
1)
      Y(61)=.0
      Y(62)=.0
      Y(122)=-2.*P*(7.*U(122)-8.*U(121)+U(120))/(X(61)-X(59))**2+
1FI(U(61)-U(122))
      DO 20 I=2,60,2
      Y(I)=-4.*EPS*P*(2.*U(I)-U(I-1)-U(I+1))/(X(I+1)-X(I-1))**2-
1FI(U(I)-U(I+61))
      Y(I+61)=-4.*P*(2.*U(I+61)-U(I+60)-U(I+62))/(X(I+1)-X(I-1))
1**2+FI(U(I)-U(I+61))
20 CONTINUE
      DO 30 I=3,59,2
      Y(I)=-2.*EPS*P*((7.*U(I)-8.*U(I-1)+U(I-2))/((X(I+2)-X(I-2))*
1(X(I)-X(I-2)))+(7.*U(I)-8.*U(I+1)+U(I+2))/((X(I+2)-X(I-2))*
2(X(I+2)-X(I))))-FI(U(I)-U(I+61))
      Y(I+61)=-2.*P*((7.*U(I+61)-8.*U(I+60)+U(I+59))/((X(I+2)-X(I-2))*
1(X(I)-X(I-2)))+(7.*U(I+61)-8.*U(I+62)+U(I+63))/((X(I+2)-X(I-2))*
2(X(I+2)-X(I))))+FI(U(I)-U(I+61))
30 CONTINUE
      RETURN
      END

```

```

C      PROGRAM TEST2(OUTPUT,TAPE6=OUTPUT)
C
C TEST PROGRAM 2 FOR J.G.VERWER'S ALGORITHM M3RK, AN EXPLICIT TIME-
C INTEGRATOR FOR SEMI-DISCRETE PARABOLIC EQUATIONS.
C THE TEST PROBLEM IS THE SINGLE TWO-DIMENSIONAL PARABOLIC EQUATION
C WHICH IS DISCUSSED IN SECTION 5.2 OF J.G.VERWER'S COMPANION PAPER
C "AN IMPLEMENTATION OF A CLASS OF STABILIZED, EXPLICIT METHODS FOR
C THE TIME INTEGRATION OF PARABOLIC EQUATIONS".
C THE TEST CONSISTS OF THE INTEGRATION OF THE SYSTEM OF 205 ODE'S
C WHICH IS CALLED SYSTEM II IN SECTION 5.2 OF THE COMPANION PAPER
C (EQUATIONS (5.6) WITH 5*41 GRID POINTS FOR THE GALERKIN DISCRETI-
C ZATION). IN THE PRESENT TEST THE TOLERANCE PARAMETER TOL=1.0E-4.
C
C *****
C
C ARRAY DECLARATIONS FOR M3RK. THE ARRAY U IS THE INPUT ARRAY FOR THE
C INITIAL VECTOR AND UOUT IS THE OUTPUT ARRAY FOR THE COMPUTED SOLUT-
C ION VECTOR.
C DER IS THE SUBROUTINE DEFINING THE SYSTEM OF ODE'S.
C
      DIMENSION INFO(15),U(205),U1(205),U2(205),DU(205),DU1(205),
+UOUT(205),SIGMA(2)
      EXTERNAL DER
C
C THE ARRAYS R AND Z WILL CONTAIN THE GRID POINTS USED BY THE GALERKIN
C DISCRETIZATION IN THE R-DIRECTION AND Z-DIRECTION,RESPECTIVELY.
C
      DIMENSION R(5),Z(41)
C
C ARRAY DECLARATIONS FOR THE AUXILARY SUBROUTINE EVAL WHICH GENERATES
C THE GALERKIN COEFFICIENTS OCCURRING IN THE SYTEM OF ODE'S(SEE EQUA-
C TIONS (5.6),(5.7) IN THE COMPANION PAPER). EVAL IS CALLED ONCE IN
C THE TEST PROGRAM.
C
      COMMON/ARRAY/DIAG(205),CO1(205),CO2(205),VRWZ(205),WZ(205)
C
C DECLARATION OF SOME CONSTANTS NEEDED IN DER. M IS THE NUMBER OF
C GRID POINTS IN THE R-DIRECTION, N IS THE NUMBER OF GRID POINTS
C IN THE Z-DIRECTION. REND IS THE RIGHT BOUNDARY OF THE R-INTERVAL
C
      COMMON/CONST/M,N,REND
C
C DEFINITION OF THE GRID FOR THE GALERKIN DISCRETIZATION. WE PUT
C M=5 AND N=41. THESE NUMBERS MAY BE CHANGED, PROVIDED THE LENGTH
C OF THE ARRAYS ARE ADJUSTED. THE VALUE OF N MINUS 1 MUST BE A
C MULTIPLE OF 4. THE INTEGER MN DEFINES THE NUMBER OF COMPONENTS
C OF THE SYSTEM OF ODE'S.
C

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      REND=1.E-4
      ZEND=0.15
      N=41
      M=5
      MN=M*N
      DR = REND/FLOAT(M-1)
      DO 5 I = 1,M
5     R(I) = DR*FLOAT(I-1)
      Z1=ZEND/10.
      N1=(N-1)/4
      DZR=Z1/FLOAT(N1)
      DZM=4.*DZR
      N2=N1+1
      N3=N2+1
      N4=3*N2-2
      N5=N4+1
      DO 6 I=1,N2
6     Z(I)=DZR*FLOAT(I-1)
      DO 7 I=N3,N4
7     Z(I)=Z(I-1)+DZM
      DO 8 I=N5,N
8     Z(I)=Z(I-1)+DZR
C
C PRINTING OF THE GRIDPOINTS.
C
      WRITE(6,20)
20  FORMAT(1H1,22H GRID POINTS ON R-AXIS)
      DO 22 I=1,M
      WRITE(6,21)R(I)
21  FORMAT(1H ,E14.8)
22  CONTINUE
      WRITE(6,23)
23  FORMAT(1H1,22H GRID POINTS ON Z-AXIS)
      DO 25 I=1,N
      WRITE(6,24)Z(I)
24  FORMAT(1H ,E14.8)
25  CONTINUE
      WRITE(6,26)
26  FORMAT(1H1)
C
C DEFINITION AND PRINTING OF THE INITIAL VECTOR OF THE SYSTEM OF
C ODE'S-UPRINT IS AN AUXILIARY PRINT ROUTINE-AND CALL OF EVAL.
C
      DO 30 I=1,M
      DO 30 J=1,N
      IJM=I+(J-1)*M
      U(IJM)=500.
30  CONTINUE
      CALL UPRINT(M,N,MN,U,Z)
      CALL EVAL(R,M,Z,N,MN)
C
C DEFINITION OF PARAMETERS FOR M3RK. T IS THE INDEPENDENT VARIABLE.
C THE MEANING OF THE OTHER PARAMETERS IS CLEARLY EXPLAINED IN THE
C PROLOGUE OF COMMENTS OF M3RK.
C
      T=0.
      TOL=1.E-4
      INFO(1)=0
      INFO(2)=2
      INFO(3)=15000
C
C INITIAL AND SUBSEQUENT CALLS OF M3RK (SEE SECTION 5.2 OF THE COM-
C PANION PAPER). TE DEFINES THE POINT WHERE OUTPUT IS DESIRED.
C
      DO 90 I=1,7
      TE=0.1
      IF(I.EQ.2)TE=.2
      IF(I.EQ.3)TE=.4
      IF(I.EQ.4)TE=.6
      IF(I.EQ.5)TE=.8
      IF(I.EQ.6)TE=.9
      IF(I.EQ.7)TE=1.0
      CALL M3RK(T,TE,MN,H,HMIN,SIGMA,TOL,
      LDER,U,U1,U2,UOUT,DU,DU1,IFLAG,INFO)
C
C PRINTING OF THE OUTPUT POINT, OF THE ERROR FLAG, OF THE ESTIMATED

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C SPECTRAL RADIUS, OF THE ARRAY INFO, AND OF THE COMPUTED SOLUTION.
C
  WRITE(6,50)TE,IFLAG,SIGMA(1)
50  FORMAT(1H1,13H ENDPOINT  =,E10.4/11H IFLAG = ,I2/
      18H SIGMA=E13.7)
  WRITE(6,60)(INFO(KK),KK=1,15)
60  FORMAT(1H ,7H INFO= ,15(15,1H ))
  CALL UPRINT(M,N,MN,UOUT,Z)
90  CONTINUE
  STOP
  END

      SUBROUTINE EVAL(R,M,Z,N,MN)
C
C THIS SUBROUTINE COMPUTES, IN A STRAIGHT FORWARD MANNER, THE STANDARD
C INTEGRALS GIVEN IN FORMULA (5.7) OF THE COMPANION PAPER. THE ARRAYS
C IN WHICH THE RESULTING VALUES ARE STORED, ARE USED IN DER. BY INSPEC-
C TION OF DER AND FORMULAS (5.6)-(5.7), THE DEFINITION OF THE ARRAYS
C WILL BECOME CLEAR.
C
  DIMENSION R(M),Z(N)
  COMMON/ARRAY/DIAG(205),CO1(205),CO2(205),VRWZ(205),WZ(205)
  WZ(1)=.0
  DO 10 I=1,MN
  DIAG(I)=.0
  CO1(I)=.0
  CO2(I)=.0
  VRWZ(I)=.0
10  CONTINUE
  M1=M-1
  N1=N-1
  DO 40 I=1,M1
  RI=R(I)
  RIPLUS=R(I+1)
  DR=RIPLUS-RI
  WLEFT=DR*(2.*RI+RIPLUS)/6.
  WRGHT=DR*(RI+2.*RIPLUS)/6.
  WZ(I)=WZ(I)+WLEFT
  WZ(I+1)=WRGHT
15  DO 30 J=1,N1
  ZJ=Z(J)
  ZJPLUS=Z(J+1)
  DZ=ZJPLUS-ZJ
  ALFA1=WLEFT/DZ
  ALFA2=WRGHT/DZ
  BETA=(WLEFT+WRGHT)*DZ/(2.*DR*DR)
  VLEFT=WLEFT*DZ/2.
  VRGHT=WRGHT*DZ/2.
  L1=M*(J-1)+I
  L2=L1+1
  L3=L1+M
  L4=L3+1
  DIAG(L1)=DIAG(L1)+ALFA1+BETA
  DIAG(L2)=DIAG(L2)+ALFA2+BETA
  DIAG(L3)=DIAG(L3)+ALFA1+BETA
  DIAG(L4)=DIAG(L4)+ALFA2+BETA
  VPWZ(L1)=VRWZ(L1)+VLEFT
  VRWZ(L3)=VRWZ(L3)+VLEFT
  VRWZ(L2)=VRWZ(L2)+VRGHT
  VPWZ(L4)=VRWZ(L4)+VRGHT
  CO1(L1)=CO1(L1)-BETA
  CO2(L1)=CO2(L1)-ALFA1
  CO2(L2)=CO2(L2)-ALFA2
  CO1(L3)=CO1(L3)-BETA
30  CONTINUE
40  CONTINUE
  WZ1=WZ(1)
  DO 50 J=1,N
  JM1=(J-1)*M+1
50  WZ(J)=VRWZ(JM1)/WZ1
  RETURN
  END

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SUBROUTINE DER(MN,P)
C
C THE PRESENT SUBROUTINE DEFINES THE SYSTEM OF ODE'S WE ARE INTEGRATING.
C THE ARRAYS DIAG, CO1, CO2, VRWZ AND WZ ARE DEFINED IN EVAL. THE ARRAYS
C P AND V ARE WORK ARRAYS.
C
  DIMENSION P(MN),V(205)
  COMMON/ARRAY/DIAG(205),CO1(205),CO2(205),VRWZ(205),WZ(205)
  COMMON/CONST/M,N,REND
C
  RLAM(X)=418.4*(((.1287496E-13*X-.116873E-9)*X+.384891E-6)
1*X-.569812E-3)*X+.57185303)
  RHOC(X)=(146.44+.018513*(X-1040.))*19300.
  BRON(X)=7.1486**2*1.E-6*(-.378808E-1+.267688E-3*X+
1.1724588E-7*X*X)/(3.141593**2*1.E-16)
  EPST(X)=(((-.270491E-17*X+.3438691E-13)*X-.163905E-9)*
1X+.3305647E-6)*X-.1194E-3)*X+.02667112
C
  DO 10 I=1,MN
10 V(I)=P(I)
  DO 20 I=1,M
  P(I)=.0
  MNI=MN-M+I
  P(MNI)=.0
20 CONTINUE
  M1=M+1
  MN1=MN-M
  DO 30 I=M1,MN1
  IMM=I-M
  IPM=I+M
  P(I)=-DIAG(I)*V(I)-CO1(I-1)*V(I-1)-CO1(I)*V(I+1)
1-CO2(IMM)*V(IMM)-CO2(I)*V(IPM)
30 CONTINUE
  N1=N-1
  DO 40 J=2,N1
  JTM=J*M
40 P(JTM)=P(JTM)-REND*WZ(J)*EPST(V(JTM))/RLAM(V(JTM))
1*5.6696E-8*V(JTM)**4
  DO 50 I=M1,MN1
50 P(I)=(P(I)*RLAM(V(I))/VRWZ(I)+BRON(V(I)))/RHOC(V(I))
  RETURN
  END

SUBROUTINE UPRINT(M,N,MN,U,Z)
C
C A PRINT ROUTINE FOR THE COMPUTED SOLUTION VECTOR.
C
  DIMENSION U(MN),Z(N)
  WRITE(6,1)
1 FORMAT(///4X,1HZ,40X,8HU(T,R,Z))
  DO 10 I=1,N
  K=(I-1)*M
  K1=K+1
  KM=K+M
10 WRITE(6,20)Z(I),(U(J),J=K1,KM)
20 FORMAT(F10.6,5X,5F13.6)
  RETURN
  END

SUBROUTINE M3RK(X,XE,N,H,HMIN,SIGMA,TOL,F,Y,
+
+ Y1,Y2,YXE,DY,DY1,IFLAG,INFO)
C*****
C* ABSTRACT *
C*****
C* M3RK IS DESIGNED TO SOLVE INITIAL VALUE PROBLEMS FOR SYSTEMS OF *
C* ORDINARY DIFFERENTIAL EQUATIONS OF THE FORM *
C* DY/DX=F(Y(1),...,Y(N)), *
C* Y(I) GIVEN AT X, *
C* WHICH ORIGINATE FROM SEMI-DISCRETIZATION OF INITIAL-BOUNDARY VALUE *
C* PROBLEMS FOR PARABOLIC PARTIAL DIFFERENTIAL EQUATIONS. M3RK IS *
C* BASED ON STABILIZED,EXPLICIT THREE-STEP RUNGE-KUTTA FORMULAS OF *
C* ORDER ONE AND TWO,OF WHICH THE DEGREE CAN VARY BETWEEN 2 AND 12. *
C*

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C* M3RK NEEDS 6 ARRAYS OF LENGTH N,WHICH ALL APPEAR IN THE CALL LIST. * 150
C* THE CODE INTEGRATES FROM X TO XE,ON NORMAL RETURN THE PARAMETERS IN * 160
C* THE CALL LIST ARE READY FOR CONTINUING THE INTEGRATION.TO CONTINUE * 170
C* THE INTEGRATION,THE USER NEEDS ONLY TO REDEFINE THE OUTPUT POINT XE * 180
C* AND CALL AGAIN. * 190
C* * 200
C* M3RK CALLS 11 SUBROUTINES WHICH HAVE BEEN WRITTEN TO STRUCTURE THE * 210
C* PROGRAM.THESE SUBROUTINES ARE: * 220
C* HSTART - HSTART COMPUTES THE INITIAL STEPLENGTH * 230
C* PARAM - PARAM COMPUTES PARAMETERS OF THE VARIOUS IMPLEMENTED * 240
C* SCHEMES FROM THE COEFFICIENTS OF STABILITY POLYNOMIALS * 250
C* POWERM - POWERM ESTIMATES THE SPECTRAL RADIUS OF THE JACOBIAN OF F * 260
C* MAXDEG - MAXDEG COMPUTES THE MAXIMAL DEGREE OF THE FORMULAS, * 270
C* WHICH IS ALLOWED WITH RESPECT TO INTERNAL STABILITY * 280
C* MINDEG - MINDEG COMPUTES THE MINIMAL DEGREE OF THE FORMULAS, * 290
C* WHICH IS ALLOWED WITH RESPECT TO ABSOLUTE STABILITY * 300
C* STEP - STEP CONTAINS THE ACTUAL INTEGRATOR * 310
C* ESTIMA - ESTIMA COMPUTES A LOCAL ERROR BOUND AND ESTIMATES A LOCAL * 320
C* ERROR * 330
C* NEWH - NEWH DELIVERS A NEW STEPLENGTH * 340
C* INTER1 - INTER1 PERFORMS INTERPOLATION AFTER A CHANGE OF THE * 350
C* STEPLENGTH * 360
C* INTER2 - INTER2 INTERPOLATES THE SOLUTION AT THE OUTPUT POINT XE * 370
C* SHIFT - SHIFT SHIFTS THE DATA FOR A NEXT STEP * 380
C* * 390
C* THESE SUBROUTINES ARE COMPLETELY LOCAL,I.E.THE INFORMATION THEY * 400
C* NEED IS PASSED THROUGH THE PARAMETER LISTS.THE WHOLE PACKAGE HAS * 410
C* BEEN TESTED ON A CDC CYBER 73-28 USING AN ARITHMETIC PRECISION OF * 420
C* 14 DIGITS.THE CODE POWERM USES THE CDC SYSTEM SUBPROGRAMS RANSET * 430
C* AND RANF CONSTITUTING A RANDOM GENERATOR.RANSET AND RANF MUST BE * 440
C* REPLACED WHEN USING THE PROGRAM ON ANOTHER COMPUTER. * 450
C* THE CODES CALLED BY M3RK USE NO MACHINE DEPENDENT CONSTANTS. * 460
C* M3RK USES ONE MACHINE DEPENDENT CONSTANT,NAMELY THE ARITHMETIC * 470
C* PRECISION OF THE COMPUTER.IN THE PROGRAM THE INTERNAL VARIABLE APR, * 480
C* WHICH REPRESENTS THE ARITHMETIC PRECISION,EQUALS 1.0E-14. APR MUST * 490
C* BE CHANGED ACCORDINGLY WHEN USING THE PROGRAM ON ANOTHER COMPUTER. * 500
C* * 510
C* THE WHOLE PROGRAM PACKAGE IS ACCEPTED BY THE PFORT VERIFIER.THE * 520
C* PFORT VERIFIER IS A PROGRAM WHICH CHECKS A FORTRAN PROGRAM,I.E. A * 530
C* MAIN PROGRAM AND SUBPROGRAMS,FOR ADHERENCE TO PFORT,A PORTABLE * 540
C* SUBSET OF AMERICAN NATIONAL STANDARD FORTRAN(SEEI1 ).THE WHOLE * 550
C* PROGRAM PACKAGE IS COMPLETELY EXPLAINED AND DESCRIBED IN I2 . * 560
C***** 570
C* MEANING OF THE PARAMETERS * 580
C***** 590
C* X - VARIABLE : INDEPENDENT VARIABLE * 600
C* XE - EXPRESSION : OUTPUT POINT AT WHICH SOLUTION IS DESIRED * 610
C* N - EXPRESSION : NUMBER OF EQUATIONS * 620
C* H - VARIABLE : STEPLENGTH * 630
C* HMIN - VARIABLE : MINIMAL STEPLENGTH * 640
C* SIGMA - ARRAY : AN ARRAY OF LENGTH 2 CONTAINING ESTIMATES * 650
C* OF THE SPECTRAL RADIUS OF THE JACOBIAN OF F * 660
C* TOL - EXPRESSION : LOCAL ERROR TOLERANCE * 670
C* F - SUBROUTINE : DERIVATIVE * 680
C* Y - ARRAY : SOLUTION VECTOR AT X,INPUT AND WORK ARRAY * 690
C* Y1 - ARRAY : SOLUTION VECTOR AT X-H,WORK ARRAY * 700
C* Y2 - ARRAY : SOLUTION VECTOR AT X-2H,WORK ARRAY * 710
C* YXE - ARRAY : SOLUTION VECTOR AT XE,OUTPUT AND WORK ARRAY * 720
C* DY - ARRAY : DERIVATIVE VECTOR AT X,WORK ARRAY * 730
C* DY1 - ARRAY : DERIVATIVE VECTOR AT X-H,WORK ARRAY * 740
C* IFLAG - VARIABLE : ERROR FLAG * 750
C* INFO - ARRAY : INTEGER ARRAY OF LENGTH 15.INFO IS USED TO * 760
C* PASS INFORMATION TO INITIALIZE THE CODE,TO * 770
C* PASS INFORMATION BETWEEN THE MAIN PROGRAM AND * 780
C* THE SUBPROGRAMS,TO PASS INFORMATION TO THE * 790
C* USER ABOUT THE STATUS OF THE INTEGRATION,AND * 800
C* TO RETAIN INFORMATION FOR SUBSEQUENT CALLS. * 810
C***** 820
C* FIRST CALL TO M3RK * 830
C***** 840
C* THE USER MUST PROVIDE STORAGE IN HIS CALLING PROGRAM FOR THE ARRAYS * 850
C* IN THE CALL LIST.HE HAS TO SUPPLY THE SUBROUTINE F(N,Y) FOR * 860
C* EVALUATING THE DERIVATIVES DY(I)/DT,I=1,...,N,WHICH MUST BE OVER- * 870
C* WRITTEN ON Y(I).FOR THE SPECTRAL RADIUS THERE EXIST THREE OPTIONS * 880
C* WHICH MUST BE SELECTED WITH INFO(2).IN INFO(3) THE USER MUST GIVE * 890

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C* A MAXIMUM FOR THE NUMBER OF EVALUATIONS OF F(Y) TO BE SPENT (ON RE- * 900
C* TURN IT MAY OCCUR THAT INFO(3) IS EXCEEDED WITH ABOUT 50) * 910
C***** * 920
C* INPUT PARAMETERS * 930
C***** * 940
C* X - INITIAL VALUE OF THE INDEPENDENT VARIABLE * 950
C* XE - OUTPUT POINT AT WHICH SOLUTION IS DESIRED * 960
C* N - NUMBER OF EQUATIONS * 970
C* SIGMA - (1)= AN UPPER ESTIMATION OF THE SPECTRAL RADIUS OF THE JA- * 980
C* COBIAN OF F IN CASE OF OPTION 1.FOR OPTION 2 AND 3 NO * 990
C* INITIALIZATION IS REQUIRED * 1000
C* TOL - LOCAL ERROR TOLERANCE * 1010
C* Y - VECTOR OF INITIAL VALUES OF THE DEPENDENT VARIABLES * 1020
C* INFO - (1)= 0 TO INDICATE FIRST CALL * 1030
C* (2)= 1 TO INDICATE OPTION 1 FOR THE SPECTRAL RADIUS,I.E. * 1040
C* THE USER MUST INITIALIZE SIGMA(1) * 1050
C* = 2 TO INDICATE OPTION 2 FOR THE SPECTRAL RADIUS,I.E. * 1060
C* THE CODE INITIALIZES SIGMA(1) AND CONTROLS SIGMA(1) * 1070
C* = 3 TO INDICATE OPTION 3 FOR THE SPECTRAL RADIUS,I.E. * 1080
C* THE CODE ONLY INITIALIZES SIGMA(1) AT THE FIRST CALL * 1090
C* (3)= MAXIMUM NUMBER OF F(Y) EVALUATIONS TO BE SPENT * 1100
C***** * 1110
C* OUTPUT PARAMETERS * 1120
C***** * 1130
C* X - LAST POINT REACHED IN INTEGRATION.NORMALLY X IS SLIGHTLY * 1140
C* BEYOND XE * 1150
C* H - INITIAL STEPLENGTH FOR SUBSEQUENT CALL * 1160
C* HMIN - MINIMAL STEPLENGTH USED BY M3RK * 1170
C* SIGMA - (1)= UPPER ESTIMATION OF THE SPECTRAL RADIUS INITIALIZED BY * 1180
C* THE USER OR BY POWERM * 1190
C* - (2)= INACCURATE ESTIMATION OF THE SPECTRAL RADIUS USED FOR * 1200
C* ITS CONTROL * 1210
C* Y - SOLUTION VECTOR AT X * 1220
C* Y1 - SOLUTION VECTOR AT X-H * 1230
C* Y2 - SOLUTION VECTOR AT X-2H * 1240
C* YXE - SOLUTION VECTOR AT OUTPUT POINT XE * 1250
C* DY - DERIVATIVE VECTOR AT X * 1260
C* DY1 - DERIVATIVE VECTOR AT X-H * 1270
C* IFLAG - = 0 NORMAL RETURN,I.E.OUTPUT POINT IS REACHED * 1280
C* = 1 OUTPUT POINT IS NOT REACHED.THE MAXIMUM NUMBER OF * 1290
C* F(Y)-EVALUATIONS HAS BEEN SPENT.THE PROCESS CAN BE CON- * 1300
C* TINUED BY INCREASING INFO(3) AND CALLING AGAIN * 1310
C* = 2 MAXIMAL DEGREE FALLS OUTSIDE THE RANGE AS TOL/APR IS * 1320
C* TOO SMALL.THE PROCESS IS NOT STARTED * 1330
C* = 3 POWERM FAILED IN THE ESTIMATION OF THE SPECTRAL RADIUS. * 1340
C* THE PROCESS IS DISCONTINUED * 1350
C* INFO - (1) = 1 TO INDICATE THAT THE NEXT CALL IS A SUBSEQUENT ONE * 1360
C* (2) = 1 IN CASE OF OPTION 1 OR 3,ELSE 2 * 1370
C* (3) = UNCHANGED * 1380
C* (4) = TOTAL NUMBER OF INTEGRATION STEPS PERFORMED,I.E. * 1390
C* ACCEPTED AND REJECTED ONES * 1400
C* (5) = NUMBER OF REJECTED INTEGRATION STEPS * 1410
C* (6) = NUMBER OF RESTARTS INITIATED BY THE CODE * 1420
C* (7) = TOTAL NUMBER OF DERIVATIVE EVALUATIONS * 1430
C* (8) = NUMBER OF DERIVATIVE EVALUATIONS USED FOR THE ESTIMA- * 1440
C* TION AND CONTROL OF THE SPECTRAL RADIUS * 1450
C* (9) = CURRENT DEGREE * 1460
C* (10)= MAXIMAL DEGREE FOR THE FIRST ORDER FORMULAS * 1470
C* (11)= MAXIMAL DEGREE FOR THE SECOND ORDER FORMULAS * 1480
C* (12)= CURRENT ORDER * 1490
C* (13)= NUMBER OF STEPS PERFORMED AFTER START OR RESTART * 1500
C* (14)= NUMBER OF STEPS PERFORMED AFTER CHANGE OF H OR ORDER * 1510
C* (15)= NUMBER OF STEPS PERFORMED AFTER ESTIMATION OF SPEC- * 1520
C* TRAL RADIUS * 1530
C* * 1540
C* IT IS EMPHASIZED THAT THE OUTPUT LIST GIVEN ABOVE IS NOT STRICTLY * 1550
C* VALID WHEN ON RETURN IFLAG IS EQUAL TO 2 OR 3. * 1560
C* IT IS FURTHER EMPHASIZED THAT ON RETURN THE ARRAY Y ALWAYS CONTAINS * 1570
C* THE SOLUTION VECTOR AT THE POINT X.THUS WHEN IFLAG=3,THE USER HAS * 1580
C* THE POSSIBILITY TO RESTART THE PROCESS AT THE POINT X,PROVIDED HE * 1590
C* IS ABLE TO TAKE ACTION WITH RESPECT TO THE ESTIMATION OF THE SPEC- * 1600
C* TRAL RADIUS. * 1610
C***** * 1620
C* SUBSEQUENT CALLS TO M3RK * 1630

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C***** 1640
C* ON RETURN OF M3RK ALL PARAMETERS ARE READY FOR CONTINUING THE INTE- * 1650
C* GRATION,PROVIDED IFLAG IS NOT EQUAL TO 2 OR 3.IF XE IS REACHED AND * 1660
C* A NORMAL CONTINUATION IS DESIRED,THE USER NEEDS ONLY TO DEFINE A NEW* 1670
C* OUTPUT POINT XE AND CALL AGAIN.IF ON RETURN IFLAG=1 AND THE USER * 1680
C* WANTS TO CONTINUE,HE ONLY NEEDS TO INCREASE INFO(3) AND CALL AGAIN. * 1690
C* THE PROGRAM IS WRITTEN IN SUCH A WAY THAT THE CHOICE OF OUTPUT * 1700
C* POINTS DOES NOT AFFECT THE INTEGRATION PROCESS ITSELF.BETWEEN SUB- * 1710
C* SEQUENT CALLS THE USER MAY INCREASE TOL.ALL OTHER PARAMETERS MUST * 1720
C* REMAIN UNCHANGED.THE COUNTERS IN INFO ARE USED ACCUMULATIVELY. * 1730
C***** 1740
C* PROGRAM TEXT * 1750
C***** 1760
      DIMENSION Y(N),Y1(N),Y2(N),YXE(N),DY(N),DY1(N),SIGMA(2),INFO(15) * 1770
C***** * 1780
C* MEANING OF THE INTERNAL VARIABLES * 1790
C***** * 1800
C* ALFA - THE FACTOR FOR CHANGING THE STEPLENGTH * 1810
C* APR - THE ARITHMETIC PRECISION * 1820
C* HOLD - PREVIOUSLY ACCEPTED STEPLENGTH * 1830
C* MOLD - PREVIOUSLY USED DEGREE * 1840
C* REJECT - NUMBER OF SUCCESSIVE STEP FAILURES * 1850
C* B - NONZERO B-PARAMETERS OF THE SCHEME * 1860
C* C - C-PARAMETERS OF THE SCHEME * 1870
C* LA - LABDA-PARAMETERS OF THE SCHEME * 1880
C* HMAX - MAXIMAL STEPSIZES WITH RESPECT TO ABSOLUTE STABILITY * 1890
C* EPS - LOCAL ERROR BOUND * 1900
C* ERROR - ESTIMATED LOCAL ERROR * 1910
C***** 1920
      INTEGER REJECT * 1930
      REAL LA * 1940
      DIMENSION B(2),C(12),LA(12),HMAX(2) * 1950
      EXTERNAL F * 1960
C***** 1970
C* IF ALREADY PAST OUTPUT POINT DURING A SUBSEQUENT CALL,THEN INTER- * 1980
C* POLATE AND RETURN * 1990
C***** 2000
      IF(INFO(1).EQ.0.OR.X.LT.XE) GOTO 10 * 2010
      CALL INTER2(N,Y,Y1,Y2,YXE,(X-XE)/H) * 2020
      RETURN * 2030
C***** 2040
C* SET THE ERROR FLAG IFLAG EQUAL TO ZERO AND INITIALIZE APR.DETERMINE * 2050
C* THE MAXIMAL DEGREES WITH RESPECT TO INTERNAL STABILITY.IF NECESSARY * 2060
C* INTERRUPT * 2070
C***** 2080
      10 IFLAG=0 * 2090
      APR=1.0E-14 * 2100
      CALL MAXDFG(TOL,APR,IFLAG,INFO) * 2110
      IF(IFLAG.NE.2) GOTO 20 * 2120
      RETURN * 2130
C***** 2140
C* SET THE CONTROL VARIABLES INFO(9),REJECT AND HOLD FOR A CONTINUING * 2150
C* CALL,AND INITIALIZE THE ARRAY HMAX * 2160
C***** 2170
      20 INFO(9)=0 * 2180
      IF(INFO(1).EQ.0) GOTO 30 * 2190
      REJECT=0 * 2200
      HOLD=H * 2210
      HMAX(1)=5.15*FLOAT(INFO(10))*FLOAT(INFO(10))/SIGMA(1) * 2220
      HMAX(2)=2.29*FLOAT(INFO(11))*FLOAT(INFO(11))/SIGMA(1) * 2230
      GOTO 80 * 2240
C***** 2250
C* SET REJECT,HOLD AND H FOR THE FIRST CALL. SET THE ELEMENTS INFO(I), * 2260
C* I=4,...,8,13,14,15 EQUAL TO ZERO AND INFO(12) EQUAL TO 2.TO PREPARE * 2270
C* THE FIRST STEP EVALUATE F(Y) AND SUBSTITUTE INTO DY.IF NECESSARY * 2280
C* ESTIMATE THE SPECTRAL RADIUS AND CHECK FOR A FAILURE OF THE POWER * 2290
C* METHOD.INITIALIZE ARRAY HMAX AND ESTIMATE THE INITIAL STEPLENGTH * 2300
C***** 2310
      30 INFO(1)=1 * 2320
      REJECT=0 * 2330
      DO 40 I=4,8 * 2340
      40 INFO(I)=0 * 2350
      DO 50 I=13,15 * 2360
      50 INFO(I)=0 * 2370
      INFO(12)=2 * 2380
      DO 60 I=1,N * 2390

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        DY(I)=Y(I)                                2400
        DY1(I)=0.                                2410
        Y1(I)=0.                                  2420
        Y2(I)=0.                                  2430
600 CONTINUE                                     2440
        CALL F(N,DY)                               2450
        INFO(7)=INFO(7)+1                         2460
        IF(INFO(2).EQ.1) GOTO 700                 2470
        SIGMA(1)=0.                                2480
        CALL POWERM(N,Y,Y1,YXE,DY,DY1,F,SIGMA,APR,IFLAG,INFO) 2490
        IF(IFLAG.NE.3) GOTO 700                   2500
        RETURN                                     2510
C                                                2520
700 HMAX(1)=5.15*FLOAT(INFO(10))*FLOAT(INFO(10))/SIGMA(1) 2530
    HMAX(2)=2.29*FLOAT(INFO(11))*FLOAT(INFO(11))/SIGMA(1) 2540
    CALL HSTART(N,Y,DY,YXE,F,TOL,APR,SIGMA(1),HMIN,INFO) 2550
    H=HMIN                                         2560
    HOLD=H                                         2570
C*****                                           2580
C* DETERMINE THE DEGREE, AND, IF NECESSARY, CALCULATE THE PARAMETERS OF * 2590
C* THE SCHEME TO BE USED                          * 2600
C*****                                           2610
800 MOLD=INFO(9)                                  2620
    CALL MINDEG(H,SIGMA(1),INFO)                  2630
    IF(INFO(9).EQ.MOLD) GOTO 900                  2640
    CALL PARAM(C,LA,B,INFO)                       2650
C*****                                           2660
C* CHECK IF THE MAXIMUM NUMBER OF EVALUATIONS IS REACHED.UPDATE HMIN * 2670
C* AND CALCULATE A SOLUTION AT X+H                * 2680
C*****                                           2690
900 IF(INFO(7).GE.INFO(3)) IFLAG=1               2700
    IF(IFLAG.NE.1) GOTO 1000                     2710
    RETURN                                         2720
C                                                2730
1000 IF(H.LT.HMIN) HMIN=H                         2740
    CALL STEP(N,Y,Y1,Y2,YXE,DY,DY1,H,F,C,LA,B,INFO) 2750
    INFO(13)=INFO(13)+1                          2760
    INFO(4)=INFO(4)+1                             2770
C*****                                           2780
C* IF THE PROCESS IS IN THE START PHASE,SHIFT THE DATA.CHECK FOR THE * 2790
C* CONTINUATION WITH A THREE-STEP SCHEME.IF THE OUTPUT POINT IS PASSED * 2800
C* INTERPOLATE AND RETURN                          * 2810
C*****                                           2820
    IF(INFO(13).GE.3) GOTO 1100                  2830
    CALL SHIFT(N,Y,Y1,Y2,YXE,DY,DY1,X,HOLD,F,INFO) 2840
    INFO(14)=INFO(14)+1                          2850
    INFO(15)=INFO(15)+1                          2860
    IF(INFO(13).EQ.1) GOTO 900                   2870
    INFO(9)=0.                                    2880
    IF(X.LT.XE) GOTO 800                         2890
    CALL INTER2(N,Y,Y1,Y2,YXE,(X-XE)/HOLD)       2900
    RETURN                                         2910
C*****                                           2920
C* CALCULATE THE LOCAL ERRORBOUND AND ESTIMATE THE LOCAL ERROR. * 2930
C*****                                           2940
1100 CALL ESTIMA(N,Y,Y1,Y2,YXE,TOL,EPS,ERROR,INFO) 2950
C*****                                           2960
C* IF THE ERROR IS TOO LARGE FOR THE THIRD STEP AFTER START,THEN RE- * 2970
C* START AT INITIAL POINT WITH H=H/10           * 2980
C*****                                           2990
    IF(EPS.GE.ERROR.OR.INFO(13).NE.3) GOTO 1300 3000
    INFO(6)=INFO(6)+1                             3010
    INFO(5)=INFO(5)+3                             3020
    INFO(9)=0.                                    3030
    INFO(13)=0.                                   3040
    INFO(14)=0.                                   3050
    INFO(15)=0.                                   3060
    X=X-2.*H                                       3070
    H=H/10.                                        3080
    HOLD=H                                         3090
    DO 1200 I=1,N                                  3100
        Y(I)=Y2(I)                                3110
        DY(I)=Y(I)                                3120
1200 CONTINUE                                     3130
        CALL F(N,DY)                               3140
        INFO(7)=INFO(7)+1                         3150

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      GOTO 80 3160
C***** 3170
C* IF STEP FAILED,CHECK FOR A REESTIMATION OF THE SPECTRAL RADIUS.IF * 3180
C* NECESSARY,CHECK FOR A FAILURE OF POWERM AND UPDATE HMAX * 3190
C***** 3200
      130 IF(INFO(2).EQ.1.OR.INFO(15).EQ.0.OR.EPS.GE.ERROR) GOTO 150 3210
      SIGMA(1)=0. 3220
      CALL POWERM(N,Y,Y1,YXE,DY,DY1,F,SIGMA,APR,IFLAG,INFO) 3230
      IF(IFLAG.NE.3) GOTO 140 3240
      RETURN 3250
C 3260
      140 HMAX(1)=5.15*FLOAT(INFO(10))*FLOAT(INFO(10))/SIGMA(1) 3270
      HMAX(2)=2.29*FLOAT(INFO(11))*FLOAT(INFO(11))/SIGMA(1) 3280
C***** 3290
C* CALCULATE A NEW STEPLENGTH * 3300
C***** 3310
      150 HOLD=H 3320
      CALL NEWH(EPS/ERROR,HOLD,H,ALFA,HMAX,INFO) 3330
C***** 3340
C* IF THE ORDER EQUALS 1 AND THE STEPLENGTH IS SMALLER THAN THE MAXI- * 3350
C* MAL STEPLENGTH FOR ORDER 2 RESET THE ORDER * 3360
C***** 3370
      IF(INFO(12).EQ.1.AND.H.LT.HMAX(2)) INFO(9)=0 3380
      IF(INFO(9).EQ.0) INFO(12)=2 3390
C***** 3400
C* IF STEP FAILED REJECT THE INTEGRATION STEP.CHECK FOR THREE SUCCES- * 3410
C* SIVE FAILURES,AND,IF NECESSARY,INTERPOLATE FOR THE NEW STEPLENGTH * 3420
C* ***** 3430
      IF(EPS.GE.ERROR) GOTO 170 3440
      REJECT=REJECT+1 3450
      INFO(5)=INFO(5)+1 3460
      IF(REJECT.EQ.3) GOTO 160 3470
      CALL INTER1(N,Y,Y1,Y2,DY1,F,ALFA,INFO) 3480
      GOTO 80 3490
C***** 3500
C* RESET REJECT,INFO(I),I=6,9,12,13,14 AND THE STEPLENGTH FOR A * 3510
C* RESTART * 3520
C***** 3530
      160 REJECT=0 3540
      INFO(6)=INFO(6)+1 3550
      INFO(9)=0 3560
      INFO(12)=2 3570
      INFO(13)=0 3580
      INFO(14)=0 3590
      CALL HSTART(N,Y,DY,YXE,F,TOL,APR,SIGMA(1),H,INFO) 3600
      HOLD=H 3610
      GOTO 80 3620
C***** 3630
C* FIND OUT IF THE ORDER SHOULD CHANGE FROM 2 TO 1.ON EXIT OF THIS * 3640
C* PROGRAM PART H SHOULD BE RESET TO HMAX(2).IF THE ORDER HAS TO BE * 3650
C* CHANGED FROM 2 TO 1 SET INFO(9)=0 * 3660
C***** 3670
      170 IF(INFO(14).LT.3.OR.INFO(12).EQ.1) GOTO 180 3680
      IF(HOLD.NE.HMAX(2).OR.H.NE.HMAX(2)) GOTO 180 3690
      INFO(12)=1 3700
      CALL ESTIMA(N,Y,Y1,Y2,YXE,TOL,EPS,ERROR,INFO) 3710
      CALL NEWH(EPS/ERROR,HOLD,H,ALFA,HMAX,INFO) 3720
      IF(ALFA.LE.1.) INFO(12)=2 3730
      H=HMAX(2) 3740
      INFO(14)=-1 3750
      IF(INFO(12).EQ.2) GOTO 180 3760
      INFO(9)=0 3770
C***** 3780
C* SHIFT THE DATA * 3790
C***** 3800
      180 CALL SHIFT(N,Y,Y1,Y2,YXE,DY,DY1,X,HOLD,F,INFO) 3810
      REJECT=0 3820
      INFO(14)=INFO(14)+1 3830
      INFO(15)=INFO(15)+1 3840
C***** 3850
C* CHECK FOR A REESTIMATION OF THE SPECTRAL RADIUS.IF NECESSARY,CHECK * 3860
C* FOR A FAILURE OF POWERM AND UPDATE HMAX * 3870
C***** 3880
      IF(INFO(2).EQ.1.OR.INFO(15).NE.25) GOTO 200 3890
      CALL POWERM(N,Y,Y1,YXE,DY,DY1,F,SIGMA,APR,IFLAG,INFO) 3900
      IF(IFLAG.NE.3) GOTO 190 3910

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RETURN 3920
C 3930
190 HMAX(1)=5.15*FLOAT(INFO(10))*FLOAT(INFO(10))/SIGMA(1) 3940
HMAX(2)=2.29*FLOAT(INFO(11))*FLOAT(INFO(11))/SIGMA(1) 3950
C***** 3960
C* IF X IS SMALLER THAN XE CONTINUE THE INTEGRATION. IF NECESSARY, IN- * 3970
C* TERPOLATE FOR A CHANGE OF H * 3980
C***** 3990
200 IF(X.GE.XE) GOTO 210 4000
IF(H.NE.HOLD) CALL INTER1(N,Y,Y1,Y2,DY1,F,ALFA,INFO) 4010
IF(H.NE.HOLD.OR.INFO(9).EQ.0) GOTO 80 4020
GOTO 90 4030
C***** 4040
C* INTERPOLATE AT XE AND, IF NECESSARY, INTERPOLATE FOR A CHANGE OF H * 4050
C* BEFORE RETURN * 4060
C***** 4070
210 CALL INTER2(N,Y,Y1,Y2,YXE,(X-XE)/HOLD) 4080
IF(H.NE.HOLD) CALL INTER1(N,Y,Y1,Y2,DY1,F,ALFA,INFO) 4090
RETURN 4100
END 4110

SUBROUTINE HSTART(N,Y,DY,YXE,F,TOL,APR,SIGMA,H,INFO) 4120
C***** 4130
C* SUBROUTINE HSTART CALCULATES THE INITIAL STEPLENGTH * 4140
C***** 4150
DIMENSION Y(N),DY(N),YXE(N),INFO(15) 4160
DO 10 I=1,N 4170
YXE(I)=Y(I)+DY(I)/SIGMA 4180
10 CONTINUE 4190
CALL F(N,YXE) 4200
INFO(7)=INFO(7)+1 4210
ETAT=0.0 4220
ETAE=0.0 4230
DO 20 I=1,N 4240
ETAT=ETAT+Y(I)*Y(I) 4250
E=YXE(I)-DY(I) 4260
ETAE=ETAE+E*E 4270
20 CONTINUE 4280
ETAT=TOL+TOL*SQRT(ETAT/LOAT(N)) 4290
ETAE=SQRT(ETAE/LOAT(N))/SIGMA+APR 4300
H=SQRT(ETAT/ETAE)/SIGMA/10.0 4310
C 4320
RMMAX=INFO(11) 4330
BETA=(0.03*RMMAX+0.44)*RMMAX*RMMAX 4340
IF(H.GT.BETA/SIGMA)H=BETA/SIGMA 4350
RETURN 4360
END 4370

SUBROUTINE PARAM(C,LA,B,INFO) 4380
C***** 4390
C* PARAM DETERMINES THE PARAMETERS OF THE INTEGRATION SCHEME.THESE * 4400
C* PARAMETERS ARE EXPRESSIONS DEPENDING ON THE COEFFICIENTS OF THE * 4410
C* STABILITY POLYNOMIALS,WHICH ARE STORED IN THE ARRAY D. DURING * 4420
C* THE START,I.E.INFO(13) IS 0 OR 1,THE PARAMETERS OF A ONE-STEP SCHE- * 4430
C* ME ARE DETERMINED. * 4440
C***** 4450
REAL LA 4460
DIMENSION C(12),LA(12),B(2),D(44),P(13),S(13),INFO(15) 4470
INTEGER ORDER 4480
DATA D(1),D(2),D(3),D(4),D(5),D(6),D(7),D(8),D(9),D(10), 4490
+D(11),D(12),D(13),D(14),D(15),D(16),D(17),D(18),D(19),D(20), 4500
+D(21),D(22),D(23),D(24),D(25),D(26),D(27),D(28),D(29),D(30)/ 4510
+.5454545454545454E+00,.32974222139503E+00,.15874540963720E-01, 4520
+.5454545454545454E+00,.32957779372404E+00,.18807742712872E-01, 4530
+.26931406295668E-03,.5454545454545454E+00,.32959633626966E+00, 4540
+.19843848141588E-01,.38370516974646E-03,.23214008199836E-05, 4550
+.5454545454545454E+00,.32940480780399E+00,.20289622974884E-01, 4560
+.43922728369494E-03,.38881393659393E-05,.12058633806519E-07, 4570
+.5454545454545454E+00,.32915730747582E+00,.20529934599533E-01, 4580
+.47014565070180E-03,.48729959290595E-05,.23293337843875E-07, 4590
+.41768893290961E-10,.5454545454545454E+00,.3294029056199E+00, 4600
+.20695785946957E-01,.48959305208277E-03,.55250561672575E-05/ 4610
DATA D(31),D(32),D(33),D(34),D(35),D(36),D(37),D(38),D(39), 4620
+D(40),D(41),D(42),D(43),D(44),D(45),D(46),D(47),D(48),D(49), 4630

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+D(50),D(51),D(52),D(53),D(54),D(55),D(56),D(57),D(58),D(59), 4640
 +D(60)/ 4650
 +.32042572866540E-07,.92217187087773E-10,.10431596770258E-12, 4660
 +.54545454545454E+00,.32904622047855E+00,.20769127247467E-01, 4670
 +.50144873108237E-03,.59538312498493E-05,.3841363904624E-07, 4680
 +.13735161731261E-09,.25579038177072E-12,.19355325549260E-15, 4690
 +.54545454545454E+00,.32902403825404E+00,.20835887364795E-01, 4700
 +.51019371265378E-03,.62671263848834E-05,.43273294211670E-07, 4710
 +.17557956622083E-09,.41529099936815E-12,.52977760638010E-15, 4720
 +.28162396623902E-18,.54545454545454E+00,.32900701770523E+00, 4730
 +.20847709891773E-01,.51474022625718E-03,.64655444450446E-05, 4740
 +.46694681553560E-07,.20545752812528E-09,.55985683650846E-12/ 4750
 DATA D(61),D(62),D(63),D(64),D(65),D(66),D(67),D(68),D(69), 4760
 +D(70),D(71),D(72),D(73),D(74),D(75),D(76),D(77),D(78),D(79), 4770
 +D(80),D(81),D(82),D(83),D(84),D(85),D(86),D(87),D(88),D(89), 4780
 +D(90)/ 4790
 +.92238940881933E-15,.84171480799272E-18,.32655765072494E-21, 4800
 +.54545454545454E+00,.32923047518706E+00,.20942105514450E-01, 4810
 +.52155114179973E-03,.66698403229501E-05,.49787290971999E-07, 4820
 +.23175109073032E-09,.69290153772379E-12,.13309584971259E-14, 4830
 +.15876152476718E-17,.10702794409632E-20,.31159779644428E-24, 4840
 +.54545454545454E+00,.32888824622955E+00,.20930564082556E-01, 4850
 +.52398048369937E-03,.67865318951772E-05,.51892263386503E-07, 4860
 +.25160762750795E-09,.80319349035922E-12,.17105594522426E-14, 4870
 +.24063258323529E-17,.21467888957759E-20,.11002739569540E-23, 4880
 +.24672233557657E-27,.45454545454545E+00,.39753050587769E+00/ 4890
 DATA D(91),D(92),D(93),D(94),D(95),D(96),D(97),D(98),D(99), 4900
 +D(100),D(101),D(102),D(103),D(104),D(105),D(106),D(107), 4910
 +D(108),D(109),D(110),D(111),D(112),D(113),D(114),D(115), 4920
 +D(116),D(117),D(118),D(119),D(120)/ 4930
 +.19106034236683E-01,.45454545454545E+00,.39769493354868E+00, 4940
 +.22675100922608E-01,.32438614977749E-03,.45454545454545E+00, 4950
 +.39767639100306E+00,.23921246939481E-01,.46221334545758E-03, 4960
 +.27945492539796E-05,.45454545454545E+00,.39786791946874E+00, 4970
 +.24509754044867E-01,.53071848010798E-03,.47002589400275E-05, 4980
 +.14585303356181E-07,.45454545454545E+00,.39811541979691E+00, 4990
 +.24864589588956E-01,.56998860293388E-03,.59138340603510E-05, 5000
 +.28300608926316E-07,.50812598486162E-10,.45454545454545E+00, 5010
 +.39786982171073E+00,.25000002282553E-01,.59148858437864E-03, 5020
 +.66757562996758E-05,.38720468964770E-07,.11144761163396E-09/ 5030
 DATA D(121),D(122),D(123),D(124),D(125),D(126),D(127), 5040
 +D(128),D(129),D(130),D(131),D(132),D(133),D(134),D(135), 5050
 +D(136),D(137),D(138),D(139),D(140),D(141),D(142),D(143), 5060
 +D(144),D(145),D(146),D(147),D(148),D(149),D(150)/ 5070
 +.12608153728000E-12,.45454545454545E+00,.39822650679417E+00, 5080
 +.25183525013803E-01,.60880314013113E-03,.72358433014328E-05, 5090
 +.46725317100106E-07,.16719594503501E-09,.31157527742621E-12, 5100
 +.23590370482110E-15,.45454545454545E+00,.39824868901869E+00, 5110
 +.25265819897767E-01,.61910411544859E-03,.76074711148850E-05, 5120
 +.52536344101559E-07,.21317642667569E-09,.50421345500406E-12, 5130
 +.64317688539028E-15,.34187163161474E-18,.45454545454545E+00, 5140
 +.39826570956750E+00,.25287844091655E-01,.62494425734623E-03, 5150
 +.78540296938450E-05,.56743151456408E-07,.24973869169769E-09, 5160
 +.68066567167935E-12,.11216260688042E-14,.10236802978533E-17/ 5170
 DATA D(151),D(152),D(153),D(154),D(155),D(156),D(157), 5180
 +D(158),D(159),D(160),D(161),D(162),D(163),D(164),D(165), 5190
 +D(166),D(167),D(168),D(169),D(170),D(171),D(172),D(173), 5200
 +D(174),D(175),D(176)/ 5210
 +.39720657964596E-21,.45454545454545E+00,.39804225208567E+00, 5220
 +.25341708058167E-01,.63151501325496E-03,.80807554649258E-05, 5230
 +.60354664402269E-07,.28111642092611E-09,.84106334638657E-12, 5240
 +.16167280397671E-14,.19299941789907E-17,.13021732956413E-20, 5250
 +.37944453664700E-24,.45454545454545E+00,.39838448104318E+00, 5260
 +.25416762495946E-01,.63697712076076E-03,.82551983508228E-05, 5270
 +.63149729888770E-07,.30629858708888E-09,.97808681916440E-12, 5280
 +.20836563130596E-14,.29320679574331E-17,.26166497017204E-20, 5290
 +.13415328284784E-23,.30092796196223E-27/ 5300
 DATA D(177),D(178),D(179),D(180),D(181),D(182),D(183),D(184), 5310
 +D(185),D(186),D(187),D(188),D(189),D(190),D(191),D(192),D(193), 5320
 +D(194),D(195),D(196),D(197),D(198),D(199),D(200),D(201),D(202), 5330
 +D(203),D(204),D(205),D(206)/ 5340
 +.58064516129032E+00,-.21559395174672E+00,-.30544696579492E-01, 5350
 +.58064516129032E+00,-.19115223031554E+00,-.29473834786102E-01, 5360
 +.10066347258135E-02,-.58064516129032E+00,-.18233307297138E+00, 5370
 +.28236544473632E-01,-.12030039601232E-02,-.15208008334815E-04, 5380
 +.58064516129032E+00,-.17833069941496E+00,-.28475246894827E-01, 5390

+- .14606302030482E-02, -.29964111364600E-04, -.21343804115613E-06, 5400
 +- .58064516129032E+00, -.17610367098315E+00, -.28260676645090E-01, 5410
 +- .15322458810336E-02, -.36586320114212E-04, -.39867529427935E-06, 5420
 +- .16226665135199E-08, -.58064516129032E+00, -.17483390114911E+00, 5430
 +- .27741186226246E-01, -.14815956989918E-02, -.36301045393729E-04/ 5440
 DATA D(207), D(208), D(209), D(210), D(211), D(212), D(213), D(214), 5450
 +D(215), D(216), D(217), D(218), D(219), D(220), D(221), D(222), D(223), 5460
 +D(224), D(225), D(226), D(227), D(228), D(229), D(230), D(231), D(232), 5470
 +D(233), D(234), D(235), D(236)/ 5480
 +- .44685007550778E-06, -.26875584507281E-08, -.62831231083352E-11, 5490
 +- .58064516129032E+00, -.17396071288671E+00, -.28684134226593E-01, 5500
 +- .17432495919146E-02, -.50845860092993E-04, -.79137773126678E-06, 5510
 +- .67374753547729E-08, -.29589252318632E-10, -.52416294063507E-13, 5520
 +- .58064516129032E+00, -.17339878373842E+00, -.28044727272007E-01, 5530
 +- .16253143532453E-02, -.45810065724231E-04, -.71155788003856E-06, 5540
 +- .64049072748140E-08, -.33259355915939E-10, -.92392278190522E-13, 5550
 +- .10625147968957E-15, -.58064516129032E+00, -.17144173377009E+00, 5560
 +- .28015092938518E-01, -.16157222373633E-02, -.46227724829057E-04, 5570
 +- .75479222345544E-06, -.74894835017511E-08, -.45979053535837E-10/ 5580
 DATA D(237), D(238), D(239), D(240), D(241), D(242), D(243), D(244), 5590
 +D(245), D(246), D(247), D(248), D(249), D(250), D(251), D(252), D(253), 5600
 +D(254), D(255), D(256), D(257), D(258), D(259), D(260), D(261), D(262), 5610
 +D(263), D(264), D(265), D(266)/ 5620
 +- .17060137796770E-12, -.35056663086597E-15, -.30628886618191E-18, 5630
 +- .58064516129032E+00, -.17300284166294E+00, -.27328009024214E-01, 5640
 +- .15056021545051E-02, -.41367453292114E-04, -.65884289780074E-06, 5650
 +- .65536146317341E-08, -.42091025159364E-10, -.17479977056317E-12, 5660
 +- .45368406820998E-15, -.66928703208412E-18, -.42844311155752E-21, 5670
 +- .58064516129032E+00, -.17229366960984E+00, -.28800272381574E-01, 5680
 +- .18596174430420E-02, -.59976978730906E-04, -.11090566945359E-05, 5690
 +- .12745935692238E-07, -.95158694657869E-10, -.46969406476975E-12, 5700
 +- .15218064136528E-14, -.31130952108508E-17, -.36466961532119E-20, 5710
 +- .18644934368193E-23, .15806451612903E+01, .15059165323919E+01/ 5720
 DATA D(267), D(268), D(269), D(270), D(271), D(272), D(273), D(274), 5730
 +D(275), D(276), D(277), D(278), D(279), D(280), D(281), D(282), D(283), 5740
 +D(284), D(285), D(286), D(287), D(288), D(289), D(290), D(291), D(292), 5750
 +D(293), D(294), D(295), D(296)/ 5760
 +- .1697894351019E+00, .15806451612903E+01, .14814748109607E+01, 5770
 +.19316031414798E+00, .62334163398843E-02, .15806451612903E+01, 5780
 +.14726556536165E+01, .20074218117966E+00, .86336976630508E-02, 5790
 +.11558084587197E-03, .15806451612903E+01, .14686532800601E+01, 5800
 +.20498325715728E+00, .99938118863291E-02, .19922348036296E-03, 5810
 +.13919201448922E-05, .15806451612903E+01, .14664262516283E+01, 5820
 +.20699571533935E+00, .10680275405545E-01, .24933405067263E-03, 5830
 +.26855039111666E-05, .10854850713601E-07, .15806451612903E+01, 5840
 +.14651564817943E+01, .20774599475456E+00, .10971861052267E-01, 5850
 +.27524368830002E-03, .35403656934423E-05, .22575321236761E-07/ 5860
 DATA D(297), D(298), D(299), D(300), D(301), D(302), D(303), D(304), 5870
 +D(305), D(306), D(307), D(308), D(309), D(310), D(311), D(312), D(313), 5880
 +D(314), D(315), D(316), D(317), D(318), D(319), D(320), D(321), D(322), 5890
 +D(323), D(324), D(325), D(326)/ 5900
 +.56574487882585E-10, .15806451612903E+01, .14642832935319E+01, 5910
 +.20956213101730E+00, .11509566107375E-01, .31097988163828E-03, 5920
 +.45630986133302E-05, .37079641130883E-07, .15682590950194E-09, 5930
 +.26933430035588E-12, .15806451612903E+01, .14637213643836E+01, 5940
 +.20948465321101E+00, .11536416747709E-01, .31784614870548E-03, 5950
 +.49116720047009E-05, .44528196080481E-07, .23503328331393E-09, 5960
 +.66870118493809E-12, .79239438466385E-15, .15806451612903E+01, 5970
 +.14617643144152E+01, .21141206884584E+00, .11805056288024E-01, 5980
 +.33440743994479E-03, .54414235293880E-05, .53918086434111E-07, 5990
 +.33075450191903E-09, .12264057386754E-11, .25181064036724E-14/ 6000
 DATA D(327), D(328), D(329), D(330), D(331), D(332), D(333), D(334), 6010
 +D(335), D(336), D(337), D(338), D(339), D(340), D(341), D(342), D(343), 6020
 +D(344), D(345), D(346), D(347), D(348), D(349), D(350), D(351), D(352)/ 6030
 +.21977616072810E-17, .15806451612903E+01, .14633254223081E+01, 6040
 +.20916387703869E+00, .11579454550239E-01, .32857305529008E-03, 6050
 +.54458388015784E-05, .56371977519294E-07, .37545083692440E-09, 6060
 +.16091005950325E-11, .42884398776166E-14, .64665832685414E-17, 6070
 +.42148457924105E-20, .15806451612903E+01, .14626162502550E+01, 6080
 +.21134531244915E+00, .12108121568554E-01, .35894153745450E-03, 6090
 +.62664422624671E-05, .69193362616297E-07, .50195179261835E-09, 6100
 +.24253292036317E-11, .77310365461803E-14, .15613921810526E-16, 6110
 +.18102819599155E-19, .91776107476727E-23/ 6120
 DATA D(353), D(354), D(355), D(356), D(357), D(358), D(359), D(360), 6130
 +D(361), D(362), D(363), D(364), D(365), D(366), D(367), D(368), D(369), 6140
 6150

C

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+D(370),D(371),D(372),D(373),D(374),D(375),D(376),D(377),D(378), 6160
+D(379),D(380),D(381),D(382),D(383),D(384),D(385),D(386),D(387), 6170
+D(388),D(389),D(390),D(391),D(392),D(393),D(394),D(395),D(396), 6180
+D(397),D(398),D(399),D(400),D(401),D(402)/ 6190
+1.,1.,.5, 6200
+1.,1.,.5,63304085.E-09,1.,1.,.5,79027358.E-09,37089254.E-10, 6210
+1.,1.,.5,85605575.E-09,56773923.E-10,12758716.E-11,1.,1.,.5, 6220
+89018496.E-09,67947003.E-10,23143226.E-11,2898388.E-12,1.,1.,.5, 6230
+91025408.E-09,74822425.E-10,30567580.E-11,6072001.E-12, 6240
+4679323.E-14,1.,1.,.5,92308224.E-09,79335426.E-10,35849723.E-11, 6250
+8800161.E-12,11108474.E-14,5648779.E-16,1.,1.,.5,93178948.E-09, 6260
+82451157.E-10,39680345.E-11,11000301.E-12,17552367.E-14/ 6270
DATA D(403),D(404),D(405),D(406),D(407),D(408),D(409),D(410), 6280
+D(411),D(412),D(413),D(414),D(415),D(416),D(417),D(418),D(419), 6290
+D(420),D(421),D(422),D(423),D(424),D(425),D(426),D(427),D(428), 6300
+D(429),D(430),D(431),D(432),D(433),D(434),D(435),D(436),D(437), 6310
+D(438),D(439),D(440)/ 6320
+14976734.E-16,5293311.E-18,1.,1.,.5,93797465.E-09,84690176.E-10, 6330
+42524183.E-11,12746945.E-12,23355062.E-14,25642831.E-16, 6340
+15495967.E-18,3962808.E-20,1.,1.,.5,94252811.E-09,86352191.E-10, 6350
+44683802.E-11,14135170.E-12,28359079.E-14,36242802.E-16, 6360
+28591262.E-18,12691526.E-20,24249737.E-23,1.,1.,.5,94597848.E-09, 6370
+87619296.E-10,46357872.E-11,15246910.E-12,32599754.E-14, 6380
+46107927.E-16,42819928.E-18,25110371.E-20,84319465.E-23, 6390
+12357105.E-25/ 6400
C 6410
ORDER=INFO(12) 6420
M=INFO(9) 6430
IF(INFO(13).LT.2) GOTO 50 6440
M1=176*(ORDER-1)+M*(M+1)/2-3 6450
M2=88+M1 6460
MM=M+1 6470
DO 10 J=1,MM 6480
M1PJ=M1+J 6490
M2PJ=M2+J 6500
P(J)=D(M1PJ) 6510
S(J)=D(M2PJ) 6520
10 CONTINUE 6530
IF(M.GT.2) GOTO 20 6540
P(4)=0.0 6550
S(4)=0.0 6560
20 DD=1.375-.6*FLOAT(ORDER-1) 6570
E=P(2)+2.0*(-P(3)+P(4)+S(4)) 6580
C(M)=(E/DD-((DD-0.5)/DD)**2)/(2.+E) 6590
LA(M)=1.0/DD-C(M) 6600
LA(M-1)=S(3)/LA(M) 6610
C(M-1)=P(3)/LA(M) 6620
B(1)=(P(2)-C(M))/LA(M) 6630
B(2)=P(1) 6640
IF(M.GT.2) GOTO 30 6650
RETURN 6660
C 6670
30 MM=M-2 6680
DO 40 J=1,MM 6690
MP2MJ=M+2-J 6700
MP1MJ=M+1-J 6710
C(J)=P(MP2MJ)/S(MP1MJ) 6720
LA(J)=S(MP2MJ)/S(MP1MJ) 6730
40 CONTINUE 6740
RETURN 6750
C 6760
50 M1=352+M*(M+1)/2-3 6770
MM=M+1 6780
DO 60 J=1,MM 6790
M1PJ=M1+J 6800
60 S(J)=D(M1PJ) 6810
B(1)=0.0 6820
B(2)=0.0 6830
C(M)=0.0 6840
LA(M)=S(1) 6850
MM=M-1 6860
DO 70 J=1,MM 6870
MP2MJ=M+2-J 6880
MP1MJ=M+1-J 6890
LA(J)=S(MP2MJ)/S(MP1MJ) 6900
C(J)=0.0 6910

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70 CONTINUE                                6920
RETURN                                     6930
END                                         6940

SUBROUTINE POWERM(N, Y, Y1, YXE, DY, DY1, F, SIGMA, APR, IFLAG, INFO) 6950
C*****                                6960
C* IF ON ENTRY SIGMA(1)=0 POWERM COMPUTES AN ESTIMATION OF THE SPECTR- * 6970
C* AL RADIUS OF THE JACOBIAN BY MEANS OF AN ADJUSTED POWER METHOD.THE * 6980
C* ITERATION IS STOPPED IF TWO SUCCESSIVE ITERATES DIFFER RELATIVELY * 6990
C* LESS THAN 0.001.THE MINIMAL NUMBER OF ITERATIONS IS 5.IF THE COMPU- * 7000
C* TATION DID NOT SUCCEED WITHIN 50 ITERATIONS AN ERRORMESSAGE IS GIV- * 7010
C* EN.AS A SAFETY MARGIN,THE LAST ITERATE IS ENLARGED WITH 10 PERCENT. * 7020
C* THE RESULT IS STORED IN SIGMA(1). * 7030
C* * 7040
C* IF ON ENTRY SIGMA(1) NOT EQUALS 0 POWERM PERFORMS THREE ITERATIONS * 7050
C* WITH THE ADJUSTED POWER METHOD.IF THE THIRD ITERATE IS MORE THAN 10 * 7060
C* PERCENT SMALLER THAN SIGMA(2),THE ITERATION IS CONTINUED AS IF * 7070
C* SIGMA(1)=0.IN THIS CASE THE THIRD ITERATE IS STORED IN SIGMA(2). * 7080
C* * 7090
C* THE POWER METHOD REQUIRES THREE WORK ARRAYS.WE USE YXE,DY AND DY1, * 7100
C* WHERE DY AND DY1 ARE OVERWRITTEN.ON ENTRY OF THIS ROUTINE DY AND * 7110
C* DY1 CONTAIN F(Y) AT Y=Y(X) AND Y=Y(X-H),RESPECTIVELY.THUS ON EXIT * 7120
C* OF THIS ROUTINE TWO EXTRA EVALUATIONS OF F ARE NECESSARY FOR RESTO- * 7130
C* RING THE DERIVATIVES. * 7140
C* * 7150
C* THIS CODE USES THE CDC SYSTEM SUBPROGRAMS RANSET AND RANF,GENE- * 7160
C* RATING RANDOM VALUES FROM THE INTERVAL I0,1 . THE ARGUMENT OF RANF * 7170
C* IS DUMMY AND IGNORED.RANSET INITIALIZES THE GENERATIVE VALUE OF RANF* 7180
C*****                                7190
DIMENSION Y(N), Y1(N), YXE(N), DY(N), DY1(N), SIGMA(2), INFO(15) 7200
REAL NORM, NORM0 7210
IF(INFO(2).EQ.3) INFO(2)=1 7220
INFO(15)=0 7230
TOLLIP=1.E+4*APR 7240
SIGM=0.0 7250
S0=0.0 7260
CALL RANSET(0) 7270
DO 10 I=1,N 7280
RA=2.*RANF(IDUM)-1. 7290
YXE(I)=DY(I) 7300
IF(Y(I).EQ.0.0) DY(I)=RA*TOLLIP 7310
IF(Y(I).NE.0.0) DY(I)=Y(I)*(1.0+TOLLIP*RA) 7320
DY1(I)=DY(I) 7330
S0=S0+DY(I)*DY(I) 7340
10 CONTINUE 7350
NORM0=TOLLIP*SQRT(S0) 7360
IF(NORM0.LT.TOLLIP)NORM0=TOLLIP 7370
CALL F(N,DY1) 7380
INFO(7)=INFO(7)+1 7390
INFO(8)=INFO(8)+1 7400
C 7410
DO 70 K=1,51 7420
IF(K.LT.51) GOTO 20 7430
IFLAG=3 7440
RETURN 7450
C 7460
20 S0=0 7470
DO 30 I=1,N 7480
S1=YXE(I)-DY1(I) 7490
S0=S0+S1*S1 7500
30 CONTINUE 7510
NORM=SQRT(S0) 7520
SIGM1=SIGM 7530
SIGM=NORM/NORM0 7540
C 7550
IF(K.EQ.3.AND.SIGMA(1).EQ.0.0) SIGMA(2)=SIGM 7560
IF(K.LE.2.OR.SIGMA(1).EQ.0.0) GOTO 40 7570
IF(SIGM.GE.0.9*SIGMA(2)) GOTO 80 7580
SIGMA(2)=SIGM 7590
SIGMA(1)=0. 7600
C 7610
40 IF(ABS(SIGM1-SIGM)/SIGM.GT.0.001.OR.K.LE.4) GOTO 50 7620
SIGMA(1)=1.1*SIGM 7630
GOTO 80 7640
50 DO 60 I=1,N 7650

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60 YXE(I)=DY(I)+(YXE(I)-DY1(I))/SIGM      7660
   CALL F(N,YXE)                          7670
   INFO(7)=INFO(7)+1                      7680
   INFO(8)=INFO(8)+1                      7690
70 CONTINUE                               7700
C                                          7710
80 DO 90 I=1,N                            7720
90 DY(I)=Y(I)                             7730
   CALL F(N,DY)                           7740
   INFO(7)=INFO(7)+1                      7750
   INFO(8)=INFO(8)+1                      7760
   IF(INFO(13).NE.0) GOTO 100            7770
   RETURN                                  7780
100 DO 110 I=1,N                          7790
110 DY1(I)=Y1(I)                         7800
   CALL F(N,DY1)                         7810
   INFO(7)=INFO(7)+1                    7820
   INFO(8)=INFO(8)+1                    7830
   RETURN                                  7840
END                                        7850

SUBROUTINE MAXDEG(TOL, APR, IFLAG, INFO)   7860
C*****                                7870
C* IN MAXDEG THE MAXIMAL DEGREES WITH RESPECT TO THE INTERNAL STABILI- * 7880
C* TY CONDITION ARE COMPUTED. IF ONE OF THESE MAXIMAL DEGREES FALLS * 7890
C* OUTSIDE THE RANGE, AN ERRORMESSAGE IS GIVEN. * 7900
C*****                                7910
   DIMENSION Q(11), INFO(15)             7920
   DATA Q(1), Q(2), Q(3), Q(4), Q(5), Q(6), Q(7), Q(8), Q(9), Q(10), Q(11) / 7930
   +3.E1, 1.E2, 7.E2, 4.E3, 3.E4, 2.E5, 9.E5, 5.E6, 3.E7, 2.E8, 1.E9 / 7940
   E=TOL/APR                             7950
   DO 10 I=2, 12                          7960
   J=13-I                                  7970
   IF(Q(J).LE.E) GOTO 20                 7980
10 CONTINUE                               7990
   IFLAG=2                                8000
   RETURN                                  8010
C                                          8020
20 INFO(10)=14-I                          8030
   DO 30 I=2, 12                          8040
   J=13-I                                  8050
   IF(100.0*Q(J).LE.E) GOTO 40          8060
30 CONTINUE                               8070
   IFLAG=2                                8080
   RETURN                                  8090
C                                          8100
40 INFO(11)=14-I                          8110
   RETURN                                  8120
END                                        8130

SUBROUTINE MINDEG(H, SIGMA, INFO)         8140
C*****                                8150
C* MINDEG DETERMINES THE MINIMAL DEGREE M WHICH STILL GIVES RISE TO A * 8160
C* STABLE INTEGRATION STEP. * 8170
C*****                                8180
   DIMENSION INFO(15)                   8190
   LOGICAL START                         8200
   L=INFO(12)+9                          8210
   MMAX=INFO(L)                          8220
   START=INFO(13).LT.2                   8230
   IF(.NOT.START) BETAC=5.15+2.86*FLOAT(1-INFO(12)) 8240
   DO 10 M=2, MMAX                       8250
   RM=M                                   8260
   IF(START) BETAC=0.03*RM+0.44          8270
   IF(H.LE.BETAC*RM*RM/SIGMA) GOTO 20   8280
10 CONTINUE                               8290
   M=MMAX                                8300
20 INFO(9)=M                             8310
   RETURN                                  8320
END                                        8330

SUBROUTINE STEP(N, Y, Y1, Y2, YXE, DY, DY1, H, F, C, LA, B, INFO) 8340
C*****                                8350

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C* STEP CONTAINS THE ACTUAL INTEGRATOR.FOR CONVENIENCE,THE ONE-STEP * 8360
C* SCHEME IS ALSO FORMULATED AS A THREE-STEP SCHEME BY INTRODUCING ZE- * 8370
C* RO-PARAMETERS.ON EXIT OF THIS ROUTINE THE NEW CALCULATED SOLUTION * 8380
C* VECTOR IS CONTAINED IN YXE. * 8390
C***** 8400
      REAL LA 8410
      DIMENSION Y(N),Y1(N),Y2(N),YXE(N),DY(N),DY1(N), 8420
+LA(12),C(12),B(2),INFO(15) 8430
      M=INFO(9) 8440
      D=1.375-.6*FLOAT(INFO(12)-1) 8450
      IF(INFO(13).LT.2) D=1.0 8460
      DO 10 I=1,N 8470
10  YXE(I)=DY(I) 8480
      IF(M.EQ.2) GOTO 40 8490
      MM=M-2 8500
C 8510
      DO 30 J=1,MM 8520
      HC=H*C(J) 8530
      HLA=H*LA(J) 8540
      DO 20 I=1,N 8550
      YXE(I)=Y(I)+HC*DY1(I)+HLA*YXE(I) 8560
20  CONTINUE 8570
      CALL F(N,YXE) 8580
      INFO(7)=INFO(7)+1 8590
30  CONTINUE 8600
C 8610
40  BM1=B(1) 8620
      BM11=1.0-BM1 8630
      HC=H*C(M-1) 8640
      HLA=H*LA(M-1) 8650
      DO 50 I=1,N 8660
50  YXE(I)=BM11*Y(I)+BM1*Y1(I)+HC*DY1(I)+HLA*YXE(I) 8670
      D1=1.-D 8680
      BM1=B(2)*D 8690
      BM11=(1.-B(2))*D 8700
      HC=H*C(M)*D 8710
      HLA=H*LA(M)*D 8720
      CALL F(N,YXE) 8730
      INFO(7)=INFO(7)+1 8740
      DO 60 I=1,N 8750
60  YXE(I)=BM11*Y(I)+BM1*Y1(I)+D1*Y2(I)+HC*DY1(I)+HLA*YXE(I) 8760
      RETURN 8770
      END 8780

      SUBROUTINE ESTIMA(N,Y,Y1,Y2,YXE,TOL,EPS,ERROR,INFO) 8790
C***** 8800
C* ESTIMA CALCULATES THE LOCAL ERROR BOUND EPS=(1+NORM(Y))*TOL FOR THE * 8810
C* MIXED ERROR TEST AND ESTIMATES THE LOCAL ERROR ERROR. * 8820
C***** 8830
      DIMENSION Y(N),Y1(N),Y2(N),YXE(N),INFO(15),CONST(2) 8840
      INTEGER ORDER 8850
      CONST(1)=4.70 8860
      CONST(2)=0.79 8870
      ORDER=INFO(12) 8880
      EPS=0.0 8890
      ERROR=0.0 8900
      IF(ORDER.EQ.2) GOTO 20 8910
      DO 10 I=1,N 8920
      YI=Y(I) 8930
      EPS=EPS+YI*YI 8940
      E=Y1(I)-YI-YI+YXE(I) 8950
      ERROR=ERROR+E*E 8960
10  CONTINUE 8970
C 8980
      GOTO 40 8990
20  DO 30 I=1,N 9000
      EPS=EPS+Y(I)*Y(I) 9010
      E=-Y2(I)+3.0*(Y1(I)-Y(I))+YXE(I) 9020
      ERROR=ERROR+E*E 9030
30  CONTINUE 9040
C 9050
40  EPS=TOL+TOL*SQRT(EPS/FLOAT(N)) 9060
      ERROR=CONST(ORDER)*SQRT(ERROR/FLOAT(N)) 9070
      RETURN 9080
      END 9090

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SUBROUTINE NEWH(EPSEERR,HOLD,H,ALFA,HMAX,INFO) 9100
C***** 9110
C* NEWH DELIVERS A NEW STEPLENGTH AND THE FACTOR ALFA BY WHICH THE * 9120
C* STEPLENGTH IS CHANGED. EPSEERR DENOTES EPS/ERROR. * 9130
C***** 9140
DIMENSION HMAX(2),INFO(15) 9150
INTEGER ORDER 9160
ALFA=1.0 9170
IF(INFO(14).GE.3.OR.EPSEERR.LE.1.0) GOTO 10 9180
RETURN 9190
C 9200
10 ORDER=INFO(12) 9210
ALFA=EPSEERR*(1./FLOAT(ORDER+1))/(2.0-FLOAT(ORDER-1)*0.4) 9220
IF(ALFA.GT.0.9.AND.ALFA.LT.1.1) ALFA=1.0 9230
IF(ALFA.NE.1.0) GOTO 20 9240
C 9250
RETURN 9260
20 IF(ALFA.GT.3.0) ALFA=3.0 9270
IF(ALFA.LT.0.1) ALFA=0.1 9280
H=HOLD*ALFA 9290
IF(H.GT.HMAX(ORDER)) H=HMAX(ORDER) 9300
ALFA=H/HOLD 9310
RETURN 9320
END 9330

SUBROUTINE INTER1(N,Y,Y1,Y2,DY1,F,ALFA,INFO) 9340
C***** 9350
C* INTER1 COMPUTES VALUES FOR Y(X-ALFA*H) AND Y(X-2*ALFA*H) FROM Y(X), * 9360
C* Y(X-H) AND Y(X-2H) BY QUADRATIC INTERPOLATION,AND COMPUTES THE DER- * 9370
C* IVATIVE AT X-H BY CALLING F. * 9380
C***** 9390
DIMENSION Y(N),Y1(N),Y2(N),DY1(N),INFO(15) 9400
REAL NU 9410
NU=2.-ALFA 9420
C12=(NU-1.)* (NU-2.)/2. 9430
C11=NU*(2.-NU) 9440
C10=NU*(NU-1.)/2. 9450
NU=2.-2.*ALFA 9460
C22=(NU-1.)* (NU-2.)/2. 9470
C21=NU*(2.-NU) 9480
C20=NU*(NU-1.)/2. 9490
C 9500
DO 10 I=1,N 9510
CY0=Y(I) 9520
CY1=Y1(I) 9530
CY2=Y2(I) 9540
Y1(I)=C12*CY2+C11*CY1+C10*CY0 9550
DY1(I)=Y1(I) 9560
Y2(I)=C22*CY2+C21*CY1+C20*CY0 9570
10 CONTINUE 9580
CALL F(N,DY1) 9590
INFO(7)=INFO(7)+1 9600
INFO(14)=0 9610
RETURN 9620
END 9630

SUBROUTINE INTER2(N,Y,Y1,Y2,YXE,A) 9640
C***** 9650
C* INTER2 COMPUTES THE SOLUTION AT THE OUTPUT POINT XE=X-A*H BY QUAD- * 9660
C* RATIC INTERPOLATION BETWEEN Y(X),Y(X-H) AND Y(X-2H).THE RESULT IS * 9670
C* STORED IN YXE. * 9680
C***** 9690
DIMENSION Y(N),Y1(N),Y2(N),YXE(N) 9700
REAL NU 9710
NU=2.-A 9720
C12=(NU-1.)* (NU-2.)/2. 9730
C11=NU*(2.-NU) 9740
C10=NU*(NU-1.)/2. 9750
DO 10 I=1,N 9760
10 YXE(I)=C12*Y2(I)+C11*Y1(I)+C10*Y(I) 9770
RETURN 9780
END 9790

```


ALGORITHM 554

BRENTM, A Fortran Subroutine for the
Numerical Solution of Systems of Nonlinear
Equations [C5]

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Key Words and Phrases: nonlinear equations, numerical solution, Brent's method

CR Categories: 4.6, 5.15

Language: Fortran

1. DESCRIPTION

BRENTM is a subroutine designed to solve a system of n nonlinear equations in n variables by using a modification of Brent's method [1]. This modification, and numerical results for our implementation, are discussed in [2].

Our subroutine does not use any techniques which attempt to obtain global convergence, and therefore convergence is only guaranteed if the initial estimate for the solution is close enough. On the other hand, our code does seem to have a large region of convergence; convergence only occurs at a zero of the function, and if the iteration is not making satisfactory progress, then BRENTM will attempt to diagnose this situation and stop the iteration with an appropriate message.

The user is only required to provide a subroutine which calculates components of the function. Note that because of the nature of Brent's method, it is advantageous to arrange the components of the function in increasing order of nonlinearity; in particular, any linear components should appear first.

The accuracy of BRENTM is controlled by the parameters FTOL and XTOL. Convergence occurs if all residuals are at most FTOL in magnitude. Thus FTOL should be chosen with care since an appropriate value depends on the scaling of the functions. Convergence also occurs if BRENTM estimates that

$$\|x - x^*\|_{\infty} \leq \text{TOL} \|x^*\|_{\infty}$$

where x^* is a solution to the system of nonlinear equations and x is the current estimate of x^* . Note that if x^* has components of widely different magnitude,

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```

DO 10 I = 1, N
  TEMP = DFLOAT(I)*H
  X(I) = TEMP*(TEMP - 1.D0)
10 CONTINUE
C
C INITIAL MAX-NORM OF THE RESIDUALS.
C
  FNORM1 = 0.D0
DO 20 I = 1, N
  CALL BVP(N,X,FVEC,I)
  FNORM1 = DMAX1(FNORM1,DABS(FVEC(I)))
20 CONTINUE
C
  NFCALL = 0
CALL BRENT1(BVP,N,X,FVEC,TOL,INFO,WA,LWA)
NFEV = NFCALL/N
C
C FINAL MAX-NORM OF THE RESIDUALS.
C
  FNORM2 = 0.D0
DO 30 I = 1, N
  CALL BVP(N,X,FVEC,I)
  FNORM2 = DMAX1(FNORM2,DABS(FVEC(I)))
30 CONTINUE
C
  WRITE (NWRITE,1000) N, FNORM1, FNORM2, NFEV, INFO, (X(I), I=1,N)
  STOP
1000 FORMAT (5X,10H DIMENSION,I5,5X //
1 5X,34H INITIAL MAX-NORM OF THE RESIDUALS,D15.7 //
2 5X,34H FINAL MAX-NORM OF THE RESIDUALS ,D15.7 //
3 5X,33H NUMBER OF FUNCTION EVALUATIONS ,I10 //
4 5X,15H EXIT PARAMETER ,18X,I10 //
5 5X,27H FINAL APPROXIMATE SOLUTION // (5X,5D15.7))
C
C LAST CARD OF SAMPLE PROGRAM.
C
END
SUBROUTINE BVP(N,X,FVEC,IFLAG)
  INTEGER N,IFLAG
  DOUBLE PRECISION X(N),FVEC(N)
  *****
C
C SUBROUTINE BVP DEFINES THE BOUNDARY VALUE PROBLEM.
C
  *****
  INTEGER INT,NFCALL
  DOUBLE PRECISION H,TEMP,TEMP1,TEMP2
  DOUBLE PRECISION DFLOAT
  COMMON /REFNUM/ NFCALL
  DFLOAT(INT) = INT
  H = 1.D0/DFLOAT(N+1)
  TEMP = 0.5D0*(X(IFLAG) + DFLOAT(IFLAG)*H + 1.D0)**3)
  TEMP1 = 0.D0
  IF (IFLAG.NE. 1) TEMP1 = X(IFLAG-1)
  TEMP2 = 0.D0
  IF (IFLAG.NE. N) TEMP2 = X(IFLAG+1)
  FVEC(IFLAG) = 2.D0*X(IFLAG) - TEMP1 - TEMP2 + TEMP*H**2
  NFCALL = NFCALL + 1
  RETURN
C
C LAST CARD OF SUBROUTINE BVP.
C
END
SUBROUTINE BRENT1(FCN,N,X,FVEC,TOL,INFO,WA,LWA)
  INTEGER N,INFO,LWA
  DOUBLE PRECISION TOL
  DOUBLE PRECISION X(N),FVEC(N),WA(LWA)
  EXTERNAL FCN
  *****
C
C SUBROUTINE BRENT1
C
  THE PURPOSE OF THIS SUBROUTINE IS TO FIND A ZERO OF
  A SYSTEM OF N NONLINEAR EQUATIONS IN N VARIABLES BY A
  METHOD DUE TO R. BRENT. THIS IS DONE BY USING THE
  MORE GENERAL NONLINEAR EQUATION SOLVER BRENTM.

```

```

C
C THE SUBROUTINE STATEMENT IS
C
C SUBROUTINE BRENT1(FCN,N,X,FVEC,TOL,INFO,WA,LWA)
C
C WHERE
C
C FCN IS THE NAME OF THE USER-SUPPLIED SUBROUTINE WHICH
C CALCULATES COMPONENTS OF THE FUNCTION. FCN SHOULD BE
C DECLARED IN AN EXTERNAL STATEMENT IN THE USER CALLING
C PROGRAM, AND SHOULD BE WRITTEN AS FOLLOWS.
C
C SUBROUTINE FCN(N,X,FVEC,IFLAG)
C INTEGER N,IFLAG
C DOUBLE PRECISION X(N),FVEC(N)
C -----
C CALCULATE THE IFLAG-TH COMPONENT OF THE FUNCTION
C AND RETURN THIS VALUE IN FVEC(IFLAG).
C -----
C RETURN
C END
C
C THE VALUE OF IFLAG SHOULD NOT BE CHANGED BY FCN UNLESS
C THE USER WANTS TO TERMINATE EXECUTION OF BRENT1.
C IN THIS CASE SET IFLAG TO A NEGATIVE INTEGER.
C
C N IS A POSITIVE INTEGER INPUT VARIABLE SET TO THE NUMBER
C OF EQUATIONS AND VARIABLES.
C
C X IS AN ARRAY OF LENGTH N. ON INPUT IT MUST CONTAIN
C AN INITIAL ESTIMATE OF THE SOLUTION VECTOR. ON OUTPUT X
C CONTAINS THE FINAL ESTIMATE OF THE SOLUTION VECTOR.
C
C FVEC IS AN ARRAY OF LENGTH N. ON OUTPUT IT CONTAINS
C THE FINAL RESIDUALS.
C
C TOL IS A NONNEGATIVE INPUT VARIABLE. THE ALGORITHM CONVERGES
C IF EITHER ALL THE RESIDUALS ARE AT MOST TOL IN MAGNITUDE,
C OR IF THE ALGORITHM ESTIMATES THAT THE RELATIVE ERROR
C BETWEEN X AND THE SOLUTION IS AT MOST TOL.
C
C INFO IS AN INTEGER OUTPUT VARIABLE SET AS FOLLOWS. IF
C THE USER HAS TERMINATED EXECUTION, INFO WILL BE SET TO
C THE (NEGATIVE) VALUE OF IFLAG. SEE DESCRIPTION OF FCN.
C OTHERWISE
C
C INFO = 0 IMPROPER INPUT PARAMETERS.
C
C INFO = 1 ALL RESIDUALS ARE AT MOST TOL IN MAGNITUDE.
C
C INFO = 2 ALGORITHM ESTIMATES THAT THE RELATIVE ERROR
C BETWEEN X AND THE SOLUTION IS AT MOST TOL.
C
C INFO = 3 CONDITIONS FOR INFO = 1 AND INFO = 2 BOTH HOLD.
C
C INFO = 4 NUMBER OF FUNCTION EVALUATIONS HAS REACHED OR
C EXCEEDED 50*(N+3).
C
C INFO = 5 APPROXIMATE JACOBIAN MATRIX IS SINGULAR.
C
C INFO = 6 ITERATION IS NOT MAKING GOOD PROGRESS.
C
C INFO = 7 ITERATION IS DIVERGING.
C
C INFO = 8 ITERATION IS CONVERGING, BUT TOL IS TOO
C SMALL, OR THE CONVERGENCE IS VERY SLOW
C DUE TO A JACOBIAN SINGULAR NEAR THE OUTPUT
C X OR DUE TO BADLY SCALED VARIABLES.
C
C WA IS A WORK ARRAY OF LENGTH LWA.
C
C LWA IS A POSITIVE INTEGER INPUT VARIABLE NOT LESS THAN
C N*(N+3).
C
C SUBPROGRAMS REQUIRED
C

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C	THE USER WANTS TO TERMINATE EXECUTION OF BRENTM.	00000370
C	IN THIS CASE SET IFLAG TO A NEGATIVE INTEGER.	00000380
C		00000390
C	N IS A POSITIVE INTEGER INPUT VARIABLE SET TO THE NUMBER OF	00000400
C	EQUATIONS AND VARIABLES.	00000410
C		00000420
C	X IS AN ARRAY OF LENGTH N. ON INPUT IT MUST CONTAIN	00000430
C	AN ESTIMATE TO THE SOLUTION OF THE SYSTEM OF EQUATIONS.	00000440
C	ON OUTPUT X CONTAINS THE FINAL ESTIMATE TO THE SOLUTION	00000450
C	OF THE SYSTEM OF EQUATIONS.	00000460
C		00000470
C	FVEC IS AN ARRAY OF LENGTH N. ON OUTPUT IT CONTAINS	00000480
C	THE FINAL RESIDUALS.	00000490
C		00000500
C	FTOL IS A NONNEGATIVE INPUT VARIABLE. CONVERGENCE	00000510
C	OCCURS IF ALL RESIDUALS ARE AT MOST FTOL IN MAGNITUDE.	00000520
C		00000530
C	XTOL IS A NONNEGATIVE INPUT VARIABLE. CONVERGENCE	00000540
C	OCCURS IF THE RELATIVE ERROR BETWEEN TWO SUCCESSIVE	00000550
C	ITERATES IS AT MOST XTOL.	00000560
C		00000570
C	MAXFEV IS A POSITIVE INTEGER INPUT VARIABLE. TERMINATION	00000580
C	OCCURS IF THE NUMBER OF FUNCTION EVALUATIONS IS AT	00000590
C	LEAST MAXFEV BY THE END OF AN ITERATION. IN BRENTM,	00000600
C	A FUNCTION EVALUATION CORRESPONDS TO N CALLS TO FCN.	00000610
C		00000620
C	MOPT IS A POSITIVE INTEGER INPUT VARIABLE. MOPT SPECIFIES	00000630
C	THE NUMBER OF TIMES THAT THE APPROXIMATE JACOBIAN IS	00000640
C	USED DURING EACH ITERATION WHICH EMPLOYS ITERATIVE	00000650
C	REFINEMENT. IF MOPT IS 1, NO ITERATIVE REFINEMENT WILL	00000660
C	BE DONE. MAXIMUM EFFICIENCY IS USUALLY OBTAINED IF	00000670
C	MOPT MAXIMIZES $\log(k+1)/(N+2*k+1)$ FOR $k = 1, \dots, N$.	00000680
C		00000690
C	INFO IS AN INTEGER OUTPUT VARIABLE SET AS FOLLOWS. IF	00000700
C	THE USER HAS TERMINATED EXECUTION, INFO WILL BE SET TO	00000710
C	THE (NEGATIVE) VALUE OF IFLAG. SEE DESCRIPTION OF FCN.	00000720
C	OTHERWISE	00000730
C		00000740
C	INFO = 0 IMPROPER INPUT PARAMETERS.	00000750
C		00000760
C	INFO = 1 ALL RESIDUALS ARE AT MOST FTOL IN MAGNITUDE.	00000770
C		00000780
C	INFO = 2 RELATIVE ERROR BETWEEN TWO SUCCESSIVE ITERATES	00000790
C	IS AT MOST XTOL.	00000800
C		00000810
C	INFO = 3 CONDITIONS FOR INFO = 1 AND INFO = 2 BOTH HOLD.	00000820
C		00000830
C	INFO = 4 NUMBER OF FUNCTION EVALUATIONS HAS REACHED OR	00000840
C	EXCEEDED MAXFEV.	00000850
C		00000860
C	INFO = 5 APPROXIMATE JACOBIAN MATRIX IS SINGULAR.	00000870
C		00000880
C	INFO = 6 ITERATION IS NOT MAKING GOOD PROGRESS.	00000890
C		00000900
C	INFO = 7 ITERATION IS DIVERGING.	00000910
C		00000920
C	INFO = 8 ITERATION IS CONVERGING, BUT XTOL IS TOO	00000930
C	SMALL, OR THE CONVERGENCE IS VERY SLOW	00000940
C	DUE TO A JACOBIAN SINGULAR NEAR THE OUTPUT	00000950
C	X OR DUE TO BADLY SCALED VARIABLES.	00000960
C		00000970
C	NFEV IS AN INTEGER OUTPUT VARIABLE SET TO THE NUMBER OF	00000980
C	FUNCTION EVALUATIONS USED IN PRODUCING X. IN BRENTM,	00000990
C	A FUNCTION EVALUATION CORRESPONDS TO N CALLS TO FCN.	00001000
C		00001010
C	Q IS AN N BY N ARRAY. IF JAC DENOTES THE APPROXIMATE	00001020
C	JACOBIAN, THEN ON OUTPUT Q IS (A MULTIPLE OF) AN	00001030
C	ORTHOGONAL MATRIX SUCH THAT JAC*Q IS A LOWER TRIANGULAR	00001040
C	MATRIX. ONLY THE DIAGONAL ELEMENTS OF JAC*Q NEED	00001050
C	TO BE STORED, AND THESE CAN BE FOUND IN SIGMA.	00001060
C		00001070
C	LDQ IS A POSITIVE INTEGER INPUT VARIABLE NOT LESS THAN N	00001080
C	WHICH SPECIFIES THE LEADING DIMENSION OF THE ARRAY Q.	00001090
C		00001100
C	SIGMA IS A LINEAR ARRAY OF LENGTH N. ON OUTPUT SIGMA	00001110
C	CONTAINS THE DIAGONAL ELEMENTS OF THE MATRIX JAC*Q.	00001120


```

DIFIT1 = DIFIT                                00001890
FNORM = ZERO                                  00001900
C                                               00001910
C COMPUTE THE STEP H FOR THE DIVIDED DIFFERENCE WHICH 00001920
C APPROXIMATES THE K-TH ROW OF THE JACOBIAN MATRIX. 00001930
C                                               00001940
H = EPS*XNORM                                  00001950
IF (H .EQ. ZERO) H = EPS                       00001960
DO 40 J = 1, N                                  00001970
  DO 30 I = 1, N                                  00001980
    Q(I,J) = ZERO                                00001990
  30 CONTINUE                                    00002000
    Q(J,J) = H                                    00002010
    WA1(J) = X(J)                                 00002020
  40 CONTINUE                                    00002030
C                                               00002040
C ENTER A SUBITERATION.                         00002050
C                                               00002060
DO 150 K = 1, N                                  00002070
  IFLAG = K                                      00002080
  CALL FCN(N,WA1,FVEC,IFLAG)                     00002090
  FKY = FVEC(K)                                   00002100
  NFCALL = NFCALL + 1                             00002110
  NFEV = NFCALL/N                                 00002120
  IF (IFLAG .LT. 0) GO TO 230                     00002130
  FNORM = DMAX1(FNORM,DABS(FKY))                  00002140
C                                               00002150
C COMPUTE THE K-TH ROW OF THE JACOBIAN MATRIX. 00002160
C                                               00002170
DO 60 J = K, N                                  00002180
  DO 50 I = 1, N                                  00002190
    WA2(I) = WA1(I) + Q(I,J)                     00002200
  50 CONTINUE                                    00002210
    CALL FCN(N,WA2,FVEC,IFLAG)                   00002220
    FKZ = FVEC(K)                                 00002230
    NFCALL = NFCALL + 1                             00002240
    NFEV = NFCALL/N                                 00002250
    IF (IFLAG .LT. 0) GO TO 230                   00002260
    SIGMA(J) = FKZ - FKY                           00002270
  60 CONTINUE                                    00002280
    FVEC(K) = FKY                                  00002290
C                                               00002300
C COMPUTE THE HOUSEHOLDER TRANSFORMATION TO REDUCE THE K-TH ROW 00002310
C OF THE JACOBIAN MATRIX TO A MULTIPLE OF THE K-TH UNIT VECTOR. 00002320
C                                               00002330
ETA = ZERO                                       00002340
DO 70 I = K, N                                  00002350
  ETA = DMAX1(ETA,DABS(SIGMA(I)))                 00002360
  70 CONTINUE                                    00002370
  IF (ETA .EQ. ZERO) GO TO 150                    00002380
  NSING = NSING - 1                               00002390
  SKNORM = ZERO                                   00002400
  DO 80 I = K, N                                  00002410
    SIGMA(I) = SIGMA(I)/ETA                       00002420
    SKNORM = SKNORM + SIGMA(I)**2                 00002430
  80 CONTINUE                                    00002440
    SKNORM = DSQRT(SKNORM)                        00002450
    IF (SIGMA(K) .LT. ZERO) SKNORM = -SKNORM      00002460
    SIGMA(K) = SIGMA(K) + SKNORM                  00002470
C                                               00002480
C APPLY THE TRANSFORMATION AND COMPUTE THE MATRIX Q. 00002490
C                                               00002500
DO 90 I = 1, N                                  00002510
  WA2(I) = ZERO                                   00002520
  90 CONTINUE                                    00002530
  DO 110 J = K, N                                  00002540
    TEMP = SIGMA(J)                               00002550
    DO 100 I = 1, N                                00002560
      WA2(I) = WA2(I) + TEMP*Q(I,J)              00002570
    100 CONTINUE                                  00002580
  110 CONTINUE                                    00002590
  DO 130 J = K, N                                  00002600
    TEMP = SIGMA(J)/(SKNORM*SIGMA(K))            00002610
    DO 120 I = 1, N                                00002620
      Q(I,J) = Q(I,J) - TEMP*WA2(I)             00002630
    120 CONTINUE                                  00002640

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130      CONTINUE                                00002650
C                                             00002660
C      COMPUTE THE SUBITERATE.                  00002670
C                                             00002680
      SIGMA(K) = SKNORM*ETA                      00002690
      TEMP = FKY/SIGMA(K)                      00002700
      IF (H*DABS(TEMP) .GT. DELTA) TEMP = DSIGN(DELTA/H,TEMP) 00002710
      DO 140 I = 1, N                          00002720
          WA1(I) = WA1(I) + TEMP*Q(I,K)         00002730
140      CONTINUE                              00002740
150      CONTINUE                              00002750
C                                             00002760
C      COMPUTE THE NORMS OF THE ITERATE AND CORRECTION VECTOR. 00002770
C                                             00002780
      XNORM = ZERO                             00002790
      DIFIT = ZERO                             00002800
      DO 160 I = 1, N                          00002810
          XNORM = DMAX1(XNORM,DABS(WA1(I)))     00002820
          DIFIT = DMAX1(DIFIT,DABS(X(I)-WA1(I))) 00002830
          X(I) = WA1(I)                        00002840
160      CONTINUE                              00002850
C                                             00002860
C      UPDATE THE BOUND ON THE CORRECTION VECTOR. 00002870
C                                             00002880
      DELTA = DMAX1(DELTA,SCALE*XNORM)         00002890
C                                             00002900
C      DETERMINE THE PROGRESS OF THE ITERATION. 00002910
C                                             00002920
      CONV = (FNORM .LT. FNORM1 .AND. DIFIT .LT. DIFIT1 .AND. 00002930
1      NSING .EQ. 0)                          00002940
      NIER6 = NIER6 + 1                       00002950
      NIER7 = NIER7 + 1                       00002960
      NIER8 = NIER8 + 1                       00002970
      IF (CONV) NIER6 = 0                     00002980
      IF (FNORM .LT. FNORM1 .OR. DIFIT .LT. DIFIT1) NIER7 = 0 00002990
      IF (DIFIT .GT. EPS*XNORM) NIER8 = 0     00003000
C                                             00003010
C      TESTS FOR CONVERGENCE.                  00003020
C                                             00003030
      IF (FNORM .LE. FTOL) INFO = 1           00003040
      IF (DIFIT .LE. XTOL*XNORM .AND. CONV) INFO = 2 00003050
      IF (FNORM .LE. FTOL .AND. INFO .EQ. 2) INFO = 3 00003060
      IF (INFO .NE. 0) GO TO 230              00003070
C                                             00003080
C      TESTS FOR TERMINATION.                 00003090
C                                             00003100
      IF (NFEV .GE. MAXFEV) INFO = 4         00003110
      IF (NSING .EQ. N) INFO = 5             00003120
      IF (NIER6 .EQ. 5) INFO = 6            00003130
      IF (NIER7 .EQ. 3) INFO = 7            00003140
      IF (NIER8 .EQ. 4) INFO = 8            00003150
      IF (INFO .NE. 0) GO TO 230             00003160
C                                             00003170
C      ITERATIVE REFINEMENT IS USED IF THE ITERATION IS CONVERGING. 00003180
C                                             00003190
      IF (.NOT. CCNV .OR. DIFIT .GT. P05*XNORM .OR. MOPT .EQ. 1) 00003200
1      GO TO 220                              00003210
C                                             00003220
C      START ITERATIVE REFINEMENT.            00003230
C                                             00003240
      DO 210 M = 2, MOPT                      00003250
          FNORM1 = FNORM                      00003260
          FNORM = ZERO                        00003270
          DO 190 K = 1, N                    00003280
              IFLAG = K                     00003290
              CALL FCN(N,WA1,FVEC,IFLAG)     00003300
              FKY = FVEC(K)                  00003310
              NFCALL = NFCALL + 1            00003320
              NFEV = NFCALL/N                00003330
              IF (IFLAG .LT. 0) GO TO 230    00003340
              FNORM = DMAX1(FNORM,DABS(FKY)) 00003350
C                                             00003360
C      ITERATIVE REFINEMENT IS TERMINATED IF IT DOES NOT 00003370
C      GIVE A REDUCTION OF THE RESIDUALS.    00003380
C                                             00003390
C      IF (FNORM .LT. FNORM1) GO TO 170     00003400

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          FNORM = FNORM1
          GO TO 220
170      CONTINUE
          TEMP = FKY/SIGMA(K)
          DO 180 I = 1, N
              WA1(I) = WA1(I) + TEMP*Q(I,K)
180      CONTINUE
190      CONTINUE
C
C      COMPUTE THE NORMS OF THE ITERATE AND CORRECTION VECTOR.
C
          XNORM = ZERO
          DIFIT = ZERO
          DO 200 I = 1, N
              XNORM = DMAX1(XNORM,DABS(WA1(I)))
              DIFIT = DMAX1(DIFIT,DABS(X(I)-WA1(I)))
              X(I) = WA1(I)
200      CONTINUE
C
C      STOPPING CRITERIA FOR ITERATIVE REFINEMENT.
C
          IF (FNORM .LE. FTOL) INFO = 1
          IF (DIFIT .LE. XTOL*XNORM) INFO = 2
          IF (FNORM .LE. FTOL .AND. INFO .EQ. 2) INFO = 3
          IF (NFEV .GE. MAXFEV .AND. INFO .EQ. 0) INFO = 4
          IF (INFO .NE. 0) GO TO 230
210      CONTINUE
220      CONTINUE
C
C      END OF THE ITERATIVE REFINEMENT.
C
          GO TO 20
C
C      TERMINATION, EITHER NORMAL OR USER IMPOSED.
C
230      CONTINUE
          IF (IFLAG .LT. 0) INFO = IFLAG
          RETURN
C
C      LAST CARD OF SUBROUTINE BRENTM.
C
          END

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ALGORITHM 555

Chow–Yorke Algorithm for Fixed Points or Zeros of C^2 Maps [C5]

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Key Words and Phrases: fixed point, zero, nonlinear system, homotopy method, continuation method, parameterized nonlinear system, zero curve of a homotopy map, fixed points of nonlinear systems, zeros of nonlinear systems

CR Categories: 5.15

Language: Fortran

DESCRIPTION

Introduction

Homotopy methods (also known as continuation methods or methods of incremental loading) for computing zeros of nonlinear systems are well known [1, 4–11, 13]. In abstract terms a homotopy method is as follows: Let B be a Banach space and $f: B \rightarrow B$ the function whose zero is desired. Let $s: B \rightarrow B$ be a simple function with a known zero. Construct a continuous map (the homotopy) $\Phi: B \times [0, 1] \rightarrow B$ such that $\Phi(x, 0) = s(x)$ and $\Phi(x, 1) = f(x)$. Then by solving the equation $\Phi(x, \lambda) = 0$ in $B \times [0, 1]$, one attempts to move from the known zero of $s(x)$ (at $\lambda = 0$) to the unknown zero of $f(x)$ (at $\lambda = 1$). There is a considerable amount of theory concerning when this procedure will work [11], but, in general, moving from a zero of $s(x)$ to a zero of $f(x)$ may or may not be possible. Sometimes λ is treated as an independent variable [4, 11], and sometimes λ is a dependent variable with arc length or some other parameter as the independent variable [7, 9, 10, 13].

For the calculation of (Brouwer) fixed points, the supporting theory is much more satisfactory. Let B be the closed unit ball (or any compact, convex subset) in E^n , and let $f: B \rightarrow B$ be a C^2 map. Chow et al. [3] have proved the following powerful theorem: Define $\rho_a: [0, 1] \times B \rightarrow E^n$ by

$$\rho_a(\lambda, x) = \lambda(x - f(x)) + (1 - \lambda)(x - a)$$

where $f: B \rightarrow B$ is a C^2 map such that the Jacobian matrix of $x - f(x)$ is nonsingular at every fixed point of f . Then for almost all a in the interior of B , the

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set $\{(\lambda, x) \mid 0 \leq \lambda < 1, x \in B, \rho_a(\lambda, x) = 0\}$ of zeros of ρ_a consists of

- (1) a finite number of closed loops (having finite length) in $(0, 1) \times B$;
- (2) a finite number of arcs (having finite length) in $(0, 1) \times B$ with endpoints in $\{1\} \times B$;
- (3) a curve of finite length starting at $(0, a)$ and ending at $(1, \bar{x})$, where $\bar{x} \in B$ is a fixed point of f .

The curves in (1), (2), and (3) are disjoint and continuously differentiable.

In other words, with probability 1 there is a zero curve of $\rho_a(\lambda, x)$ emanating from $(0, a)$ which reaches a fixed point \bar{x} of f (at $\lambda = 1$). This zero curve is smooth and has finite length. Thus computing a fixed point of f merely amounts to tracking a (smooth) zero curve of $\rho_a(\lambda, x)$. Note that the assumptions on f (C^2 smoothness and nonsingularity of the Jacobian matrix of $x - f(x)$ at the fixed point) are quite mild considering that global convergence is guaranteed with probability 1.

A brief outline of how the zero curve is followed is given here. Complete details of the algorithm and convergence proofs are in [13]. The basic idea, which has been used by several authors [7, 9, 10], is to parameterize the zero curve of $\rho_a(\lambda, x)$ by arc length and then solve an initial value problem. The zero curve $(\lambda(s), x(s))$ emanating from $(0, a)$ satisfies

$$\rho_a(\lambda(s), x(s)) = 0, \quad \lambda(0) = 0, \quad x(0) = a \quad (1)$$

and thus is given by the solution of the initial value problem

$$\begin{aligned} \frac{d}{ds} \rho_a(\lambda(s), x(s)) &= 0, \\ \left\| \left(\frac{d\lambda}{ds}, \frac{dx}{ds} \right) \right\| &= 1, \quad \lambda(0) = 0, \quad x(0) = a. \end{aligned} \quad (2)$$

Computing $(d\lambda/ds, dx/ds)$ reduces to finding the kernel of the $n \times (n+1)$ matrix $[a - f(x), I - \lambda Df(x)]$ which always has full rank [3, 13]. ($Df(x)$ is the Jacobian matrix of $f(x)$.) Therefore there are no "singular points" along the curve.

Of course the zero curve of $\rho_a(\lambda, x)$ could be followed by a general curve tracking program like that of Kubicek [10], but since the ultimate objective is a fixed point and not the zero curve of ρ_a , special tactics are called for. There are several other reasons for not using the program in [10]. That program produces a specified number of data points on the curve, and it would be difficult to make it stop exactly at $\lambda(s) = 1$. Furthermore, the numerical linear algebra and ordinary differential equation (ODE) techniques in that program are primitive compared with current technology (as, for example, [2] and [12]).

Computing Zeros. This algorithm for computing fixed points can be used to find zeros also, but then global convergence is not guaranteed. The homotopy map for zeros is

$$\Phi_a(\lambda, x) = \lambda f(x) + (1 - \lambda)(x - a),$$

and the rest of the details are similar to the fixed point case. Observe that in the fixed point case, the zero curve of ρ_a has finite length (in fact, fairly short for practical problems), and thus stiffness of the ODE is never a problem. However, the zero curve of ϕ_a can (and in practice frequently does) wander off to infinity, and the ODE can be stiff.

There are numerous conditions on f guaranteeing that the zero curve of Φ_a reaches a zero [3, 13], but these are frequently not satisfied for practical problems. One such sufficient condition is that for some $r > 0$,

$$x'f(x) \geq 0 \quad \text{on} \quad \|x\| = r.$$

Then $f(x)$ has a zero in the ball $\|x\| \leq r$, and for almost all a in the interior of this ball, there is a zero curve of Φ_a connecting $(0, a)$ to $(1, \bar{x})$, and the Jacobian matrix $D\Phi_a$ has full rank along this zero curve [13].

In view of the preceding comments, the computer code should be used with caution on zero finding problems.

Organizational Details

There are seven subroutines: FIXPT, FODE, DCPOSE, INNER, STEP, INTRP, and ROOT. The user need only call FIXPT; all the others are directly or indirectly used by FIXPT.

The subroutines STEP, INTRP, and ROOT are from [12]. STEP is used to solve the initial value problem, and INTRP and ROOT are used to calculate the vector $x(s)$ corresponding to $\lambda(s) = 1$. FODE specifies the ordinary differential equation for STEP. FODE must determine the kernel of a matrix, and it uses DCPOSE and INNER for this purpose. DCPOSE was taken from [2] with minor modifications.

FIXPT and FODE contain dimension statements which limit N , and also use labeled common. These restrictions could have been avoided by considerably lengthening the call lists of STEP, FODE, and FIXPT. Instead we chose to use STEP verbatim, so that the user needs only one version of STEP, and future improvements in STEP can be easily incorporated into this fixed point package. The difficulty arises because FIXPT has parameters which must be passed to FODE via STEP, and DCPOSE must return information back up through FODE and STEP to FIXPT.

FIXPT limits the number of steps to 1000, but this can be changed in the DATA statement that sets the value of LIMITD.

FIXPT and FODE contain dimension statements which limit N to 100.

STEP and ROOT use machine dependent constants, for which appropriate DATA statements must be chosen, as explained in the listing. No other modifications are required by the user.

The user must supply two subroutines, F(X, V) and FJAC (X, V, K). Subroutine F evaluates f at X and returns the result in (the vector) V. Subroutine FJAC evaluates the Kth column of the Jacobian matrix of f at X and returns the result in V. FJAC may, of course, produce finite difference approximations to the Jacobian matrix. The effect of finite difference approximations on the overall efficiency and accuracy of the algorithm has not been explored so far.

Modifications for Large Sparse Problems. FIXPT was designed for low-dimensional ($n \leq 100$) problems where the Jacobian matrix of f is dense. It is not difficult to modify FIXPT for a problem where n is very large but the Jacobian matrix of f is sparse. Write the Jacobian matrix of the homotopy map as

$$D\rho_a(x, \lambda) = [I - \lambda Df(x), a - f(x)].$$

Note that the order of the variables has been switched. The subroutine DCPOSE essentially reduces $D\rho_a(x, \lambda)$ to upper triangular form. If $Df(x)$ is sparse, then $D\rho_a(x, \lambda)$ can be reduced to upper triangular form efficiently by, e.g., plane rotations. Probably DCPOSE should be tailored to the particular structure of $Df(x)$. FODE computes the kernel of the upper triangular matrix produced by DCPOSE. This calculation in FODE would have to be changed and again should probably be tailored to the particular form of $Df(x)$. Thus FODE and DCPOSE would require major changes, but all the other subroutines would remain the same (except, of course, for the DIMENSION statements in FIXPT). It might also be desirable to reduce the storage requirements of STEP as explained in [12].

Testing. FIXPT has been tested on several hundred problems of mixed polynomial, exponential, and trigonometric types, with n ranging from 2 to 100. Some problems, for example, an exponential with a rapidly oscillating trigonometric exponent, give the method a great deal of difficulty. On the other hand, high-degree (≥ 10) polynomials in 100 dimensions are handled easily. When used properly, the performance of FIXPT has been entirely satisfactory. The computed

Table I

n	ND	Arc length	FIXPT time	ZSYSTEM time
5	52	2.71	0.5	0.5
10	74	3.73	2.3	17.3*
15	97	4.49	6.9	†
20	73	5.16	10.3	†
25	81	5.70	19.6	†
30	108	6.19	40.3	†
35	121	6.65	69.1	†
40	121	7.08	98.0	†
45	123	7.48	134.7	†
50	129	7.85	187.8	†

* Did not converge.

† No convergence, reported "singular Jacobian."

fixed points are usually accurate to within the tolerance EPS and have never been in error by more than 10 EPS. Some typical computational results are given in [13]. Shown in Table I are the results for Brown's function

$$f_1(x) = x_1 - \left(\prod_{i=1}^n x_i - 1 \right),$$

$$f_j(x) = x_j - \left(\sum_{i=1}^n x_i + x_j - (n+1) \right), \quad j = 2, \dots, n,$$

which was suggested by R. Saigal as a particularly difficult problem (because the Jacobian matrix is ill conditioned). ZSYSTEM is from the IMSL library and uses a quasi-Newton method. ND is the number of Jacobian evaluations, and the CPU time is execution time (in seconds) on a CDC 6500. EPS = 1.E-8, ARCTOL = 1.0E-3 for N = 5, 10, 20, 25, and ARCTOL = 1.0E-5 for N = 15, 30, 35, 40, 45, 50. For this problem the zero curve does not turn back, although in the next example from [13] the zero curve turns back many times. Therefore, the use of the parameter λ as a *dependent* variable is crucial.

$$f_k(x) = \exp\left(\cos\left(k \sum_{i=1}^5 x_i\right)\right), \quad k = 1, \dots, 5.$$

Starting from zero, both Steffensen's method and ZSYSTEM failed to find a fixed point of $f(x)$. FIXPT required 6.39 seconds of CPU time, 513 Jacobian evaluations, with an arc length of 14.83, to compute a fixed point accurate to eight places.

Referee's Note. The single precision program given in the listing was run successfully in double precision on an IBM 3033, using the AUTODBL feature of the Fortran Extended H Compiler. However, the author reports that this worked on another system only after the final argument (0.0) in the two references to INNER in DCPOSE was changed to double precision.

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ALGORITHM

```

          DIMENSION Y(101)                                30
C                                                    40
C                                                    50
C MAIN PROGRAM AND USER WRITTEN SUBROUTINES FOR THE FUNCTION 60
C F WHOSE KTH COMPONENT IS                               70
C   EXP(COS(K*(X(1)+X(2)+...+X(N))))                     80
C                                                    90
C                                                    100
C N=3                                                    110
C ARCTOL=1.0E-4                                         120
C EPS=1.0E-8                                            130
C CALL SECOND(TIME1)                                    140
C IFLAG=0                                              150
C CALL FIXPT(N,Y,ARCTOL,EPS,ARCLN,NFE,IFLAG)           160
C NP1=N+1                                              170
C CALL SECOND(TIME2)                                    180
C WRITE (6,30) ARCLN,IFLAG,(Y(L),L=1,NP1)              190
30  FORMAT(/11H ARC LENGTH,1PE18.9,4X,4HCODE,I3/(1X,7E17.9))
C WRITE (6,40) NFE                                       210
40  FORMAT(/18,21H JACOBIAN EVALUATIONS/)
C EXTIME=TIME2-TIME1                                    230
C WRITE (6,50) EXTIME                                    240
50  FORMAT(22H EXECUTION TIME(SEC) =,F8.2/1H1)
60  CONTINUE                                           260
C STOP                                                 270
C END                                                  280

          SUBROUTINE SECOND(T)
C
C ***** THIS IS A DUMMY ROUTINE *****
C
C T SHOULD BE SET EQUAL TO A COUNTER COUNTING SECONDS
C
C   T = 0.0
C   RETURN
C   END

          SUBROUTINE F(X,V)                                290
C
C
C SUBROUTINE TO EVALUATE THE FUNCTION WHOSE JTH COMPONENT IS
C   EXP(COS(J*(X(1)+...+X(N))))
C
C ON RETURN V CONTAINS F(X)
C
C
C REAL X(100),V(100)                                    300
C COMMON /FIXEDP/ YPOLD(101),A(100),NFE,N,NP1,IFLAG    310
C SUMX=0.0                                              320
C DO 20 J=1,N                                           330
20  SUMX=SUMX+X(J)                                       340
C DO 30 K=1,N                                           350
30  V(K)=EXP(COS(FLOAT(K)*SUMX))
C RETURN                                               370
C END                                                  380

          SUBROUTINE FJAC(X,V,K)                          390
C
C

```

```

C      EVALUATES THE JACOBIAN MATRIX OF THE FUNCTION WHOSE JTH
C      COMPONENT IS
C      EXP(COS(J*(X(1)+...+X(N))))
C
C      ON RETURN V CONTAINS THE KTH COLUMN OF THE JACOBIAN MATRIX
C      DF(X) .
C
C
C      REAL X(100),V(100)                                400
C      COMMON /FIXEDP/ YPOLD(101),A(100),NFE,N,NP1,IFLAG  410
C      IF (K .GT. 1) RETURN                               420
C      SUMX=0.0                                           430
C      DO 30 J=1,N                                       440
30    SUMX=SUMX+X(J)                                     450
40    DO 60 J=1,N                                       460
C      FJ=FLOAT(J)
C      ECKX=EXP(COS(FJ*SUMX))
60    V(J)=-FJ*SIN(FJ*SUMX)*ECKX
C      RETURN                                             490
C
C
C      END OF TEST PROGRAMS.                             500
C
C
C      *****                                           510
C
C      *****                                           520
C
C      *****                                           530
C
C      *****                                           540
C
C      *****                                           550
C
C      *****                                           560
C
C
C      SUBROUTINE FIXPT(N,Y,ARCTOL,EPS,ARCLN,NFE,IFLAG)   10
C
C      SUBROUTINE FIXPT FINDS A FIXED POINT OR ZERO OF THE  20
C      N-DIMENSIONAL VECTOR FUNCTION F(X). FOR THE FIXED POINT  30
C      PROBLEM F(X) IS ASSUMED TO BE A C2 MAP OF SOME BALL INTO  40
C      ITSELF. THE EQUATION X = F(X) IS SOLVED BY          50
C      FOLLOWING THE ZERO CURVE OF THE HOMOTOPY MAP        60
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A) ,        70
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A) ,        80
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A) ,        90
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A) ,       100
C
C      STARTING FROM LAMBDA = 0, X = A. THE CURVE IS PARAMETERIZED  110
C      BY ARC LENGTH S, AND IS FOLLOWED BY SOLVING THE ORDINARY  120
C      DIFFERENTIAL EQUATION D(HOMOTOPY MAP)/DS = 0 FOR    130
C      Y(S) = (LAMBDA(S), X(S)).                            140
C
C
C      FOR THE ZERO FINDING PROBLEM F(X) IS ASSUMED TO BE A C2 MAP  150
C      SUCH THAT FOR SOME R .GT. 0, X*F(X) .GE. 0 WHENEVER NORM(X) = R.  160
C      THE EQUATION F(X) = 0 IS SOLVED BY FOLLOWING THE ZERO CURVE  170
C      OF THE HOMOTOPY MAP                                  180
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          190
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          200
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          210
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          220
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          230
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          240
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          250
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          260
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          270
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          280
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          290
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          300
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          310
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          320
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          330
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          340
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          350
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          360
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          370
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          380
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          390
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          400
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          410
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          420
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          430
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          440
C
C      LAMBDA*(X - F(X)) + (1 - LAMBDA)*(X - A)          450

```

C		460
C	Y IS AN ARRRAY OF LENGTH N + 1. (Y(2),...,Y(N+1)) = A IS THE	470
C	STARTING POINT FOR THE ZERO CURVE.	480
C		490
C	ARCTOL IS THE LOCAL ERROR ALLOWED THE ODE SOLVER WHEN	500
C	FOLLOWING THE ZERO CURVE. IF ARCTOL .LE. 0.0 ON INPUT	510
C	IT IS RESET TO .5*SQRT(EPS). NORMALLY ARCTOL SHOULD	520
C	BE CONSIDERABLY LARGER THAN EPS.	530
C		540
C	EPS IS THE LOCAL ERROR ALLOWED THE ODE SOLVER WHEN VERY	550
C	NEAR THE FIXED POINT(ZERO). EPS IS APPROXIMATELY THE	560
C	ABSOLUTE ERROR IN THE COMPUTED FIXED POINT(ZERO).	570
C		580
C	IFLAG CAN BE -1, 0, 2, OR 3. IFLAG SHOULD BE 0 ON THE FIRST	590
C	CALL TO FIXPT FOR THE PROBLEM X=F(X), AND -1 FOR THE	600
C	PROBLEM F(X)=0. IN CERTAIN SITUATIONS IFLAG IS SET TO	610
C	2 OR 3 BY FIXPT, AND FIXPT CAN BE CALLED AGAIN WITHOUT	620
C	CHANGING IFLAG.	630
C		640
C	Y, ARCTOL, EPS, ARCLN, NFE, AND IFLAG SHOULD ALL BE	650
C	VARIABLES IN THE CALLING PROGRAM.	660
C		670
C		680
C	ON OUTPUT:	690
C		700
C	N IS UNCHANGED.	710
C		720
C	Y(1) = LAMBDA, (Y(2),...,Y(N+1)) = X, AND Y IS AN APPROXIMATE	730
C	ZERO OF THE HOMOTOPY MAP. NORMALLY LAMBDA = 1 AND X IS A	740
C	FIXED POINT(ZERO) OF F(X). IN ABNORMAL SITUATIONS LAMBDA	750
C	MAY ONLY BE NEAR 1 AND X IS NEAR A FIXED POINT(ZERO).	760
C		770
C	ARCTOL = EPS AFTER A NORMAL RETURN (IFLAG = 1).	780
C		790
C	EPS IS UNCHANGED AFTER A NORMAL RETURN (IFLAG = 1). IT IS	800
C	INCREASED TO AN APPROPRIATE VALUE ON THE RETURN IFLAG = 2.	810
C		820
C	ARCLN IS THE LENGTH OF THE PATH FOLLOWED.	830
C		840
C	NFE IS THE NUMBER OF FUNCTION EVALUATIONS (= NUMBER OF	850
C	JACOBIAN EVALUATIONS).	860
C		870
C	IFLAG =	880
C	-1 CAUSES FIXPT TO INITIALIZE EVERYTHING FOR THE PROBLEM	890
C	F(X) = 0 (USE ON FIRST CALL).	900
C		910
C	0 CAUSES FIXPT TO INITIALIZE EVERYTHING FOR THE PROBLEM	920
C	X = F(X) (USE ON FIRST CALL).	930
C		940
C	1 NORMAL RETURN.	950
C		960
C	2 SPECIFIED ERROR TOLERANCE CANNOT BE MET. EPS HAS BEEN	970
C	INCREASED TO A SUITABLE VALUE. TO CONTINUE, JUST CALL	980
C	FIXPT AGAIN WITHOUT CHANGING ANY PARAMETERS.	990
C		1000
C	3 STEP HAS BEEN CALLED 1000 TIMES. TO CONTINUE, CALL	1010
C	FIXPT AGAIN WITHOUT CHANGING ANY PARAMETERS.	1020
C		1030
C	4 JACOBIAN MATRIX DOES NOT HAVE FULL RANK. THE ALGORITHM	1040
C	HAS FAILED (THE ZERO CURVE OF THE HOMOTOPY MAP CANNOT BE	1050
C	FOLLOWED ANY FURTHER).	1060
C		1070
C	5 EPS (OR ARCTOL) IS TOO LARGE. THE PROBLEM SHOULD BE	1080
C	RESTARTED BY CALLING FIXPT WITH A SMALLER EPS (OR	1090
C	ARCTOL) AND IFLAG = 0 (-1).	1100
C		1110
C	6 I - DF(X) IS NEARLY SINGULAR AT THE FIXED POINT(DF(X) IS	1120
C	NEARLY SINGULAR AT THE ZERO). ANSWER MAY NOT BE	1130
C	ACCURATE.	1140
C		1150
C	7 ILLEGAL INPUT PARAMETERS, A FATAL ERROR.	1160


```

DO 70 JW=1,N
AOLD=A(JW)
IF (IFLAGC .LT. 0)
+ A(JW)=Y(1)*YP(JW)/(1.0 - Y(1)) + Y(JW+1)
IF (IFLAGC .EQ. 0)
+ A(JW)=(Y(JW+1) - Y(1)*YP(JW))/(1.0 - Y(1))
IF (ABS(A(JW)-AOLD) .GT. .95) GO TO 40
70 CONTINUE
GO TO 100
80 IF (Y(1) .LE. .99 .OR. ST99) GO TO 100
C WHEN LAMBDA REACHES .99, THE PROBLEM WILL BE RESTARTED WITH
C A NEW A VECTOR.
90 ST99=.TRUE.
EPSSTP=EPS
ARCTOL=EPS
GO TO 60

C
C TAKE A STEP ALONG THE CURVE.
100 CALL STEP(S,Y,FODE,NP1,H,EPSSTP,WT,START,HOLD,K,KOLD,CRASH,
+ PHI,P,YP,PSI)
NFE=NFEC
C CHECK IF THE STEP WAS SUCCESSFUL.
IF (IFLAGC .NE. 4) GO TO 120
IFLAG=4
RETURN
120 IF (.NOT. CRASH) GO TO 130
C RETURN CODE FOR ERROR TOLERANCE TOO SMALL.
IFLAG=2
C CHANGE ERROR TOLERANCES.
EPS=EPSSTP
IF (ARCTOL .LT. EPS) ARCTOL=EPS
C CHANGE LIMIT ON NUMBER OF ITERATIONS.
LIMS=LIMIT-ITER
GO TO 220

C
130 EPSSTP=ARCTOL
IF (ABS(YP(1)) .LE. CURTOL) EPSSTP=EPS
IF (Y(1) .LT. 1.0) GO TO 150
IF (ST99) GO TO 160

C
C IF LAMBDA .GE. 1.0 BUT THE PROBLEM HAS NOT BEEN RESTARTED
C WITH A NEW A VECTOR, BACK UP AND RESTART.
C
S99=S-.5*HOLD
C GET AN APPROXIMATE ZERO Y(S) WITH Y(1)=LAMBDA .LT. 1.0
135 CALL INTRP(S,Y,S99,WT,P,NP1,KOLD,PHI,PSI)
IF (WT(1) .LT. 1.0) GO TO 140
S99=.5*(S-HOLD+S99)
GO TO 135

C
140 DO 144 JUDY=1,NP1
Y(JUDY)=WT(JUDY)
YPOLD(JUDY)=P(JUDY)
144 WT(JUDY)=1.0
S=S99
GO TO 90

C
150 CONTINUE
C
C ***** END OF MAIN LOOP. *****
C
C LAMBDA HAS NOT REACHED 1 IN 1000 STEPS.
IFLAG=3
RETURN

C
C
C USE INVERSE INTERPOLATION TO GET THE ANSWER AT LAMBDA = 1.0 .
C
160 SA=S-HOLD
SB=S
LCODE=1
170 CALL ROOT(SOUT,Y1SOUT,SA,SB,EPS,EPS,LCODE)
C ROOT FINDS S SUCH THAT Y(1)(S) = LAMBDA = 1 .
IF (LCODE .GT. 0) GO TO 190
CALL INTRP(S,Y,SOUT,WT,P,NP1,KOLD,PHI,PSI)
Y1SOUT=WT(1)-1.0

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```

        GO TO 170
190  IFLAG=1
C SET IFLAG = 6 IF ROOT COULD NOT GET LAMBDA = 1.0
    IF (LCODE .GT. 2) IFLAG=6
    ARCLEN=ARCLEN+SA
C LAMBDA(SA) = 1.0
    CALL INTRP(S,Y,SA,WT,P,NP1,KOLD,PHI,PSI)
C
    DO 210 J=1,NP1
210  Y(J)=WT(J)
    RETURN
220  LIMIT=LIMS
    RETURN
    END
2800

        SUBROUTINE FODE(S,Y,YP)
C
C SUBROUTINE FODE IS USED BY SUBROUTINE STEP TO SPECIFY THE
C ORDINARY DIFFERENTIAL EQUATION DY/DS = G(S,Y) , WHOSE SOLUTION
C IS THE ZERO CURVE OF THE HOMOTOPY MAP. S = ARC LENGTH,
C YP = DY/DS, AND Y(S) = (LAMBDA(S), X(S)) .
C
C ***** ARRAY DECLARATIONS. *****
C
C THE FOLLOWING ARRAYS ARE FOR A MAXIMUM OF N = 100 VARIABLES.
C TO HANDLE UP TO N VARIABLES, CHANGE EACH 100 AND 101 TO N
C AND N + 1 RESPECTIVELY.
C
    REAL Y(1),YP(1)
C ARRAYS FOR COMPUTING THE JACOBIAN MATRIX AND ITS KERNAL.
    REAL QR(100,101),ALPHA(100),TZ(101)
    INTEGER PIVOT(101)
C
    REAL INNER
    COMMON /FIXEDP/ YPOLD(101),A(100),NFE,N,NP1,IFLAG
    NDIM=100
C
C ***** END OF DIMENSIONAL INFORMATION. *****
C
    NFE=NFE+1
C NFE CONTAINS THE NUMBER OF JACOBIAN EVALUATIONS.
C * * * * *
C
C COMPUTE THE JACOBIAN MATRIX, STORE IT IN QR.
C
C    QR = ( A - F(X), I - LAMBDA*DF(X) ) .
C
C
    CALL F(Y(2),TZ)
    IF (IFLAG .LT. 0) GO TO 140
    DO 100 J=1,N
100  QR(J,1)=A(J)-TZ(J)
    DO 120 K=1,N
        CALL FJAC(Y(2),TZ,K)
        KP1=K+1
    DO 110 J=1,N
110  QR(J,KP1)=-Y(1)*TZ(J)
120  QR(K,KP1)=1.0+QR(K,KP1)
    GO TO 10
C
C    QR = ( F(X) - X + A, LAMBDA*DF(X) + (1 - LAMBDA)*I ) .
C
C
140  DO 150 J=1,N
150  QR(J,1)=TZ(J)-Y(J+1)+A(J)
    DO 170 K=1,N
        CALL FJAC(Y(2),TZ,K)
        KP1=K+1
    DO 160 J=1,N
160  QR(J,KP1)=Y(1)*TZ(J)
170  QR(K,KP1)=1.0-Y(1)+QR(K,KP1)
C
C * * * * *
C REDUCE THE JACOBIAN MATRIX TO UPPER TRIANGULAR FORM.
10  CALL DCPOSE(NDIM,N,QR,ALPHA,PIVOT,IERR,TZ,YP)
    IF (IERR .EQ. 0) GO TO 20
    IFLAG=4
2400

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RETURN 3410
C COMPUTE KERNAL OF JACOBIAN, WHICH SPECIFIES YP=DY/DS. 3420
20 TZ(NP1)=1.0 3430
DO 40 LW=1,N 3440
    I=NP1-LW 3450
    IK=I+1 3460
    SUM=0.0 3470
    DO 30 J=IK,NP1 3480
30 SUM=SUM+QR(I,J)*TZ(J) 3490
40 TZ(I)=-SUM/ALPHA(I) 3500
    YPNORM=SQRT(INNER(1,NP1,TZ,TZ,0.0)) 3510
    DO 60 K=1,NP1 3520
        KPIV=PIVOT(K)
60 YP(KPIV)=TZ(K)/YPNORM 3530
    IF (INNER(1,NP1,YP,YPOLD,0.0) .GE. 0.0) GO TO 80 3540
    DO 70 I=1,NP1 3550
70 YP(I)=-YP(I) 3560
C 3570
C SAVE CURRENT DERIVATIVE (= TANGENT VECTOR) IN YPOLD 3580
80 DO 90 I=1,NP1 3590
90 YPOLD(I)=YP(I) 3600
    RETURN 3610
    END 3620

REAL FUNCTION INNER(K,M,A,B,C) 3770
DIMENSION A(1),B(1) 3780
SUM=C 3790
DO 10 I=K,M 3800
10 SUM=SUM+A(I)*B(I) 3810
INNER=SUM 3820
RETURN 3830
END 3840

SUBROUTINE DCPOSE(NDIM,N,QR,ALPHA,PIVOT,IERR,Y,SUM) 3850
C 3860
C SUBROUTINE DCPOSE IS A MODIFICATION OF THE ALGOL PROCEDURE 3870
C DECOMPOSE IN P. BUSINGER AND G. H. GOLUB, LINEAR LEAST 3880
C SQUARES SOLUTIONS BY HOUSEHOLDER TRANSFORMATIONS, 3890
C NUMER. MATH. 7 (1965) 269-276. 3900
C 3910
C INTEGER NDIM,N,PIVOT(1) 3920
REAL QR(NDIM,1),ALPHA(N) 3930
INTEGER IERR,I,J,JBAR,K 3940
REAL BETA,SIGMA,ALPHAK,QRKK,Y(1),SUM(1) 3950
REAL INNER 3960
IERR=0 3970
NP1=N+1 3980
DO 20 J=1,NP1 3990
    SUM(J)=INNER(1,N,QR(1,J),QR(1,J),0.0) 4000
20 PIVOT(J)=J 4010
DO 500 K=1,N 4020
    SIGMA=SUM(K) 4030
    JBAR=K 4040
    KPI=K+1 4050
    DO 40 J=KPI,NP1 4060
        IF (SIGMA .GE. SUM(J)) GO TO 40 4070
        SIGMA=SUM(J) 4080
        JBAR=J 4090
40 CONTINUE 4100
    IF (JBAR .EQ. K) GO TO 70 4110
    I=PIVOT(K) 4120
    PIVOT(K)=PIVOT(JBAR) 4130
    PIVOT(JBAR)=I 4140
    SUM(JBAR)=SUM(K) 4150
    SUM(K)=SIGMA 4160
    DO 50 I=1,N 4170
        SIGMA=QR(I,K) 4180
        QR(I,K)=QR(I,JBAR) 4190
        QR(I,JBAR)=SIGMA 4200
50 CONTINUE 4210

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C      END OF COLUMN INTERCHANGE. 4220
70    SIGMA=INNER(K,N,QR(1,K),QR(1,K),0.0) 4230
      IF (SIGMA .NE. 0.0) GO TO 60 4240
      IERR=1 4250
      RETURN 4260
60    IF (K .EQ. N) GO TO 500 4270
      QRKK=QR(K,K) 4280
      ALPHAK=-SQRT(SIGMA) 4290
      IF (QRKK .LT. 0.0) ALPHAK=-ALPHAK 4300
      ALPHA(K)=ALPHAK 4310
      BETA=1.0/(SIGMA-QRKK*ALPHAK) 4320
      QR(K,K)=QRKK-ALPHAK 4330
      DO 80 J=K+1,NP1 4340
80    Y(J)=BETA*INNER(K,N,QR(1,K),QR(1,J),0.0) 4350
      DO 100 J=K+1,NP1 4360
        DO 90 I=K,N 4370
          QR(I,J)=QR(I,J)-QR(I,K)*Y(J) 4380
90    CONTINUE 4390
      SUM(J)=SUM(J)-QR(K,J)**2 4400
100  CONTINUE 4410
500  CONTINUE 4420
      ALPHA(N)=QR(N,N) 4430
      RETURN 4440
      END 4450

      SUBROUTINE STEP(X,Y,F,NEQN,H,EPS,WT,START, 4460
1    HOLD,K,KOLD,CRASH,PHI,P,YP,PSI) 4470
C 4480
C SUBROUTINE STEP INTEGRATES A SYSTEM OF FIRST ORDER ORDINARY 4490
C DIFFERENTIAL EQUATIONS ONE STEP NORMALLY FROM X TO X+H, USING A 4500
C MODIFIED DIVIDED DIFFERENCE FORM OF THE ADAMS PECE FORMULAS. LOCAL 4510
C EXTRAPOLATION IS USED TO IMPROVE ABSOLUTE STABILITY AND ACCURACY. 4520
C THE CODE ADJUSTS ITS ORDER AND STEP SIZE TO CONTROL THE LOCAL ERROR 4530
C PER UNIT STEP IN A GENERALIZED SENSE. SPECIAL DEVICES ARE INCLUDED 4540
C TO CONTROL ROUND OFF ERROR AND TO DETECT WHEN THE USER IS REQUESTING 4550
C TOO MUCH ACCURACY. 4560
C 4570
C THIS CODE IS COMPLETELY EXPLAINED AND DOCUMENTED IN THE TEXT, 4580
C COMPUTER SOLUTION OF ORDINARY DIFFERENTIAL EQUATIONS& THE INITIAL 4590
C VALUE PROBLEM BY L. F. SHAMPINE AND M. K. GORDON. 4600
C 4610
C 4620
C THE PARAMETERS REPRESENT& 4630
C X -- INDEPENDENT VARIABLE 4640
C Y(*) -- SOLUTION VECTOR AT X 4650
C YP(*) -- DERIVATIVE OF SOLUTION VECTOR AT X AFTER SUCCESSFUL 4660
C STEP 4670
C NEQN -- NUMBER OF EQUATIONS TO BE INTEGRATED 4680
C H -- APPROPRIATE STEP SIZE FOR NEXT STEP. NORMALLY DETERMINED BY 4690
C CODE 4700
C EPS -- LOCAL ERROR TOLERANCE. MUST BE VARIABLE 4710
C WT(*) -- VECTOR OF WEIGHTS FOR ERROR CRITERION 4720
C START -- LOGICAL VARIABLE SET .TRUE. FOR FIRST STEP, .FALSE. 4730
C OTHERWISE 4740
C HOLD -- STEP SIZE USED FOR LAST SUCCESSFUL STEP 4750
C K -- APPROPRIATE ORDER FOR NEXT STEP (DETERMINED BY CODE) 4760
C KOLD -- ORDER USED FOR LAST SUCCESSFUL STEP 4770
C CRASH -- LOGICAL VARIABLE SET .TRUE. WHEN NO STEP CAN BE TAKEN, 4780
C .FALSE. OTHERWISE 4790
C THE ARRAYS PHI, PSI ARE REQUIRED FOR THE INTERPOLATION SUBROUTINE 4800
C INTRP. THE ARRAY P IS INTERNAL TO THE CODE. 4810
C 4820
C INPUT TO STEP 4830
C 4840
C FIRST CALL -- 4850
C 4860
C THE USER MUST PROVIDE STORAGE IN HIS DRIVER PROGRAM FOR ALL ARRAYS 4870
C IN THE CALL LIST, NAMELY 4880
C 4890
C DIMENSION Y(NEQN),WT(NEQN),PHI(NEQN,16),P(NEQN),YP(NEQN),PSI(12) 4900
C 4910
C THE USER MUST ALSO DECLARE START AND CRASH LOGICAL VARIABLES 4920
C AND F AN EXTERNAL SUBROUTINE, SUPPLY THE SUBROUTINE F(X,Y,YP) 4930
C TO EVALUATE 4940
C DY(I)/DX = YP(I) = F(X,Y(1),Y(2),...,Y(NEQN)) 4950

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C* DATA TWOU,FOURU /Z3C200000,Z3C400000/
C* FOR THE IBM 360/370 (DOUBLE PRECISION) -- *
C* DATA TWOU,FOURU /Z342000000000000,Z344000000000000/
C* FOR THE UNIVAC 1108 AND HONEYWELL 6000 -- *
C* DATA TWOU,FOURU /2.0**(-25),2.0**(-24)/
C* FOR THE PDP-11 -- *
C* DATA TWOU,FOURU /2.0**(-22),2.0**(-21)/
C*****5700
  DATA TWO/2.0,4.0,8.0,16.0,32.0,64.0,128.0,256.0,512.0,1024.0, 5710
  1 2048.0,4096.0,8192.0/ 5720
  DATA GSTR/0.500,0.0833,0.0417,0.0264,0.0188,0.0143,0.0114,0.00936,5730
  1 0.00789,0.00679,0.00592,0.00524,0.00468/ 5740
  DATA G(1),G(2)/1.0,0.5/,SIG(1)/1.0/ 5750
C 5760
C 5770
C *** BEGIN BLOCK 0 5780
C CHECK IF STEP SIZE OR ERROR TOLERANCE IS TOO SMALL FOR MACHINE 5790
C PRECISION. IF FIRST STEP, INITIALIZE PHI ARRAY AND ESTIMATE A 5800
C STARTING STEP SIZE. 5810
C *** 5820
C 5830
C IF STEP SIZE IS TOO SMALL, DETERMINE AN ACCEPTABLE ONE 5840
C 5850
  CRASH = .TRUE. 5860
  IF(ABS(H) .GE. FOURU*ABS(X)) GO TO 5 5870
  H = SIGN(FOURU*ABS(X),H) 5880
  RETURN 5890
  5 P5EPS = 0.5*EPS 5900
C 5910
C IF ERROR TOLERANCE IS TOO SMALL, INCREASE IT TO AN ACCEPTABLE VALUE 5920
C 5930
  ROUND = 0.0 5940
  DO 10 L = 1,NEQN 5950
10 ROUND = ROUND + (Y(L)/WT(L))**2 5960
  ROUND = TWOU*SQRT(ROUND) 5970
  IF(P5EPS .GE. ROUND) GO TO 15 5980
  EPS = 2.0*ROUND*(1.0 + FOURU) 5990
  RETURN 6000
15 CRASH = .FALSE. 6010
  IF(.NOT.START) GO TO 99 6020
C 6030
C INITIALIZE. COMPUTE APPROPRIATE STEP SIZE FOR FIRST STEP 6040
C 6050
  CALL F(X,Y,YP) 6060
  SUM = 0.0 6070
  DO 20 L = 1,NEQN 6080
  PHI(L,1) = YP(L) 6090
  PHI(L,2) = 0.0 6100
20 SUM = SUM + (YP(L)/WT(L))**2 6110
  SUM = SQRT(SUM) 6120
  ABSH = ABS(H) 6130
  IF(EPS .LT. 16.0*SUM*H) ABSH = 0.25*SQRT(EPS/SUM) 6140
  H = SIGN(AMAX1(ABSH,FOURU*ABS(X)),H) 6150
  HOLD = 0.0 6160
  K = 1 6170
  KOLD = 0 6180
  START = .FALSE. 6190
  PHASE1 = .TRUE. 6200
  NORND = .TRUE. 6210
  IF(P5EPS .GT. 100.0*ROUND) GO TO 99 6220
  NORND = .FALSE. 6230
  DO 25 L = 1,NEQN 6240
25 PHI(L,15) = 0.0 6250
99 IFAIL = 0 6260
C *** END BLOCK 0 *** 6270
C 6280
C *** BEGIN BLOCK 1 *** 6290
C COMPUTE COEFFICIENTS OF FORMULAS FOR THIS STEP. AVOID COMPUTING 6300
C THOSE QUANTITIES NOT CHANGED WHEN STEP SIZE IS NOT CHANGED. 6310
C *** 6320
C 6330
100 KP1 = K+1 6340
  KP2 = K+2 6350
  KM1 = K-1 6360
  KM2 = K-2 6370
C 6380

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C   NS IS THE NUMBER OF STEPS TAKEN WITH SIZE H, INCLUDING THE CURRENT 6390
C   ONE.  WHEN K.LT.NS, NO COEFFICIENTS CHANGE 6400
      IF(H .NE. HOLD) NS = 0 6410
      NS = MIN0(NS+1,KOLD+1) 6420
      NSP1 = NS+1 6430
      IF (K .LT. NS) GO TO 199 6440
C 6450
C   COMPUTE THOSE COMPONENTS OF ALPHA(*),BETA(*),PSI(*),SIG(*) WHICH 6460
C   ARE CHANGED 6470
C 6480
      BETA(NS) = 1.0 6490
      REALNS = NS 6500
      ALPHA(NS) = 1.0/REALNS 6510
      TEMP1 = H*REALNS 6520
      SIG(NSP1) = 1.0 6530
      IF(K .LT. NSP1) GO TO 110 6540
      DO 105 I = NSP1,K 6550
        IM1 = I-1 6560
        TEMP2 = PSI(IM1) 6570
        PSI(IM1) = TEMP1 6580
        BETA(I) = BETA(IM1)*PSI(IM1)/TEMP2 6590
        TEMP1 = TEMP2 + H 6600
        ALPHA(I) = H/TEMP1 6610
        REALI = I 6620
      105 SIG(I+1) = REALI*ALPHA(I)*SIG(I) 6630
      110 PSI(K) = TEMP1 6640
C 6650
C   COMPUTE COEFFICIENTS G(*) 6660
C 6670
C   INITIALIZE V(*) AND SET W(*).  G(2) IS SET IN DATA STATEMENT 6680
C 6690
      IF(NS .GT. 1) GO TO 120 6700
      DO 115 IQ = 1,K 6710
        TEMP3 = IQ*(IQ+1) 6720
        V(IQ) = 1.0/TEMP3 6730
      115 W(IQ) = V(IQ) 6740
      GO TO 140 6750
C 6760
C   IF ORDER WAS RAISED, UPDATE DIAGONAL PART OF V(*) 6770
C 6780
      120 IF(K .LE. KOLD) GO TO 130 6790
      TEMP4 = K*KP1 6800
      V(K) = 1.0/TEMP4 6810
      NSM2 = NS-2 6820
      IF(NSM2 .LT. 1) GO TO 130 6830
      DO 125 J = 1,NSM2 6840
        I = K-J 6850
      125 V(I) = V(I) - ALPHA(J+1)*V(I+1) 6860
C 6870
C   UPDATE V(*) AND SET W(*) 6880
C 6890
      130 LIMIT1 = KP1 - NS 6900
      TEMP5 = ALPHA(NS) 6910
      DO 135 IQ = 1,LIMIT1 6920
        V(IQ) = V(IQ) - TEMP5*V(IQ+1) 6930
      135 W(IQ) = V(IQ) 6940
      G(NSP1) = W(1) 6950
C 6960
C   COMPUTE THE G(*) IN THE WORK VECTOR W(*) 6970
C 6980
      140 NSP2 = NS + 2 6990
      IF(KP1 .LT. NSP2) GO TO 199 7000
      DO 150 I = NSP2,KP1 7010
        LIMIT2 = KP2 - I 7020
        TEMP6 = ALPHA(I-1) 7030
        DO 145 IQ = 1,LIMIT2 7040
          145 W(IQ) = W(IQ) - TEMP6*W(IQ+1) 7050
        150 G(I) = W(1) 7060
      199 CONTINUE 7070
C 7080
      *** END BLOCK 1 *** 7090
C 7100
      *** BEGIN BLOCK 2 *** 7110
C   PREDICT A SOLUTION P(*), EVALUATE DERIVATIVES USING PREDICTED 7120
C   SOLUTION, ESTIMATE LOCAL ERROR AT ORDER K AND ERRORS AT ORDERS K, 7130
C   K-1, K-2 AS IF CONSTANT STEP SIZE WERE USED. 7140
      ***

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C 7150
C CHANGE PHI TO PHI STAR 7160
C 7170
  IF(K .LT. NSP1) GO TO 215 7180
  DO 210 I = NSP1,K 7190
    TEMP1 = BETA(I) 7200
    DO 205 L = 1,NEQN 7210
      PHI(L,I) = TEMP1*PHI(L,I) 7220
    210 CONTINUE 7230
C 7240
C PREDICT SOLUTION AND DIFFERENCES 7250
C 7260
  215 DO 220 L = 1,NEQN 7270
    PHI(L,KP2) = PHI(L,KP1) 7280
    PHI(L,KP1) = 0.0 7290
  220 P(L) = 0.0 7300
    DO 230 J = 1,K 7310
      I = KP1 - J 7320
      IP1 = I+1 7330
      TEMP2 = G(I) 7340
      DO 225 L = 1,NEQN 7350
        P(L) = P(L) + TEMP2*PHI(L,I) 7360
  225 PHI(L,I) = PHI(L,I) + PHI(L,IP1) 7370
  230 CONTINUE 7380
    IF(NORND) GO TO 240 7390
    DO 235 L = 1,NEQN 7400
      TAU = H*P(L) - PHI(L,15) 7410
      P(L) = Y(L) + TAU 7420
  235 PHI(L,16) = (P(L) - Y(L)) -TAU 7430
    GO TO 250 7440
  240 DO 245 L = 1,NEQN 7450
  245 P(L) = Y(L) + H*P(L) 7460
  250 XOLD = X 7470
    X = X + H 7480
    ABSH = ABS(H) 7490
    CALL F(X,P,YP) 7500
C 7510
C ESTIMATE ERRORS AT ORDERS K,K-1,K-2 7520
  ERKM2 = 0.0 7530
  ERKM1 = 0.0 7540
  ERK = 0.0 7550
  DO 265 L = 1,NEQN 7560
    TEMP3 = 1.0/WT(L) 7570
    TEMP4 = YP(L) - PHI(L,1) 7580
    IF(KM2) 265,260,255 7590
  255 ERKM2 = ERKM2 + ((PHI(L,KM1)+TEMP4)*TEMP3)**2 7600
  260 ERKM1 = ERKM1 + ((PHI(L,K)+TEMP4)*TEMP3)**2 7610
  265 ERK = ERK + (TEMP4*TEMP3)**2 7620
    IF(KM2) 280,275,270 7630
  270 ERKM2 = ABSH*SIG(KM1)*GSTR(KM2)*SQRT(ERKM2) 7640
  275 ERKM1 = ABSH*SIG(K)*GSTR(KM1)*SQRT(ERKM1) 7650
  280 TEMP5 = ABSH*SQRT(ERK) 7660
    ERR = TEMP5*(G(K)-G(KP1)) 7670
    ERK = TEMP5*SIG(KP1)*GSTR(K) 7680
    KNEW = K 7690
C 7700
C TEST IF ORDER SHOULD BE LOWERED 7710
C 7720
  IF(KM2) 299,290,285 7730
  285 IF(AMAX1(ERKM1,ERKM2) .LE. ERK) KNEW = KM1 7740
  GO TO 299 7750
  290 IF(ERKM1 .LE. 0.5*ERK) KNEW = KM1 7760
C 7770
C TEST IF STEP SUCCESSFUL 7780
C 7790
  299 IF(ERR .LE. EPS) GO TO 400 7800
C *** END BLOCK 2 *** 7810
C 7820
C *** BEGIN BLOCK 3 *** 7830
C THE STEP IS UNSUCCESSFUL. RESTORE X, PHI(*,*), PSI(*) 7840
C IF THIRD CONSECUTIVE FAILURE, SET ORDER TO ONE. IF STEP FAILS MORE 7850
C THAN THREE TIMES, CONSIDER AN OPTIMAL STEP SIZE. DOUBLE ERROR 7860
C TOLERANCE AND RETURN IF ESTIMATED STEP SIZE IS TOO SMALL FOR MACHINE 7870
C PRECISION 7880
C *** 7890
C 7900

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C   RESTORE X, PHI(*,*) AND PSI(*)          7910
C                                           7920
C   PHASE1 = .FALSE.                       7930
C   X = XOLD                                7940
C   DO 310 I = 1,K                          7950
C     TEMP1 = 1.0/BETA(I)                   7960
C     IP1 = I+1                             7970
C     DO 305 L = 1,NEQN                     7980
305   PHI(L,I) = TEMP1*(PHI(L,I) - PHI(L,IP1)) 7990
310   CONTINUE                             8000
C     IF(K .LT. 2) GO TO 320                8010
C     DO 315 I = 2,K                       8020
315   PSI(I-1) = PSI(I) - H                8030
C                                           8040
C   ON THIRD FAILURE, SET ORDER TO ONE.    8050
C   THEREAFTER, USE OPTIMAL STEP          8060
C   SIZE                                   8070
C                                           8080
320   IFAIL = IFAIL + 1                    8090
C     TEMP2 = 0.5                          8100
C     IF(IFAIL - 3) 335,330,325            8110
325   IF(P5EPS .LT. 0.25*ERK) TEMP2 = SQRT(P5EPS/ERK) 8120
330   KNEW = 1                             8130
335   H = TEMP2*H                          8140
C     K = KNEW                             8150
C     IF(ABS(H) .GE. FOURU*ABS(X)) GO TO 340 8160
C     CRASH = .TRUE.                       8170
C     H = SIGN(FOURU*ABS(X),H)             8180
C     EPS = EPS + EPS                      8190
C     RETURN                                8200
340   GO TO 100                             8210
C     ***      END BLOCK 3      ***        8220
C                                           8230
C     ***      BEGIN BLOCK 4      ***      8240
C   THE STEP IS SUCCESSFUL. CORRECT THE    8250
C   PREDICTED SOLUTION, EVALUATE          8260
C   THE DERIVATIVES USING THE CORRECTED   8270
C   SOLUTION AND UPDATE THE               8280
C   DIFFERENCES. DETERMINE BEST ORDER    8290
C   AND STEP SIZE FOR NEXT STEP.         8300
C     ***                                     8310
400   KOLD = K                             8320
C     HOLD = H                             8330
C                                           8340
C   CORRECT AND EVALUATE                  8350
C                                           8360
C     TEMP1 = H*G(KP1)                     8370
C     IF(NORND) CO TO 410                  8380
C     DO 405 L = 1,NEQN                    8390
C       RHO = TEMP1*(YP(L) - PHI(L,1)) - PHI(L,16) 8400
C       Y(L) = P(L) + RHO                 8410
405   PHI(L,15) = (Y(L) - P(L)) - RHO      8420
C     GO TO 420                             8430
410   DO 415 L = 1,NEQN                   8440
415   Y(L) = P(L) + TEMP1*(YP(L) - PHI(L,1)) 8450
420   CALL F(X,Y,YP)                      8460
C                                           8470
C   UPDATE DIFFERENCES FOR NEXT STEP      8480
C                                           8490
C     DO 425 L = 1,NEQN                   8500
C       PHI(L,KP1) = YP(L) - PHI(L,1)     8510
425   PHI(L,KP2) = PHI(L,KP1) - PHI(L,KP2) 8520
C     DO 435 I = 1,K                       8530
C       DO 430 L = 1,NEQN                 8540
430   PHI(L,I) = PHI(L,I) + PHI(L,KP1)    8550
435   CONTINUE                             8560
C                                           8570
C   ESTIMATE ERROR AT ORDER K+1 UNLESS&   8580
C   IN FIRST PHASE WHEN ALWAYS RAISE     8590
C   ORDER,                               8600
C   ALREADY DECIDED TO LOWER ORDER,     8610
C   STEP SIZE NOT CONSTANT SO ESTIMATE   8620
C   UNRELIABLE                          8630
C                                           8640
C     ERKP1 = 0.0                          8650
C     IF(KNEW .EQ. KM1 .OR. K .EQ. 12)    8660
C     PHASE1 = .FALSE.
C     IF(PHASE1) GO TO 450
C     IF(KNEW .EQ. KM1) GO TO 455
C     IF(KP1 .GT. NS) GO TO 460
C     DO 440 L = 1,NEQN
440   ERKP1 = ERKP1 + (PHI(L,KP2)/WT(L))**2
C     ERKP1 = ABSH*GSTR(KP1)*SQRT(ERKP1)

```

```

C                                     8670
C USING ESTIMATED ERROR AT ORDER K+1, DETERMINE APPROPRIATE ORDER 8680
C FOR NEXT STEP 8690
C 8700
C IF(K .GT. 1) GO TO 445 8710
C IF(ERKP1 .GE. 0.5*ERK) GO TO 460 8720
C GO TO 450 8730
445 IF(ERKM1 .LE. AMIN1(ERK,ERKP1)) GO TO 455 8740
C IF(ERKP1 .GE. ERK .OR. K .EQ. 12) GO TO 460 8750
C 8760
C HERE ERKP1 .LT. ERK .LT. AMAX1(ERKM1,ERKM2) ELSE ORDER WOULD HAVE 8770
C BEEN LOWERED IN BLOCK 2. THUS ORDER IS TO BE RAISED 8780
C 8790
C RAISE ORDER 8800
C 8810
450 K = KP1 8820
C ERK = ERKP1 8830
C GO TO 460 8840
C 8850
C LOWER ORDER 8860
C 8870
455 K = KM1 8880
C ERK = ERKM1 8890
C 8900
C WITH NEW ORDER DETERMINE APPROPRIATE STEP SIZE FOR NEXT STEP 8910
C 8920
460 HNEW = H + H 8930
C IF(PHASE1) GO TO 465 8940
C IF(P5EPS .GE. ERK*TWO(K+1)) GO TO 465 8950
C HNEW = H 8960
C IF(P5EPS .GE. ERK) GO TO 465 8970
C TEMP2 = K+1 8980
C R = (P5EPS/ERK)**(1.0/TEMP2) 8990
C HNEW = ABSH*AMAX1(0.5,AMIN1(0.9,R)) 9000
C HNEW = SIGN(AMAX1(HNEW,FOURU*ABS(X)),H) 9010
465 H = HNEW 9020
C RETURN 9030
C *** END BLOCK 4 *** 9040
C END 9050

SUBROUTINE INTRP(X,Y,XOUT,YOUT,YPOUT,NEQN,KOLD,PHI,PSI) 9060
C 9070
C THE METHODS IN SUBROUTINE STEP APPROXIMATE THE SOLUTION NEAR X 9080
C BY A POLYNOMIAL. SUBROUTINE INTRP APPROXIMATES THE SOLUTION AT 9090
C XOUT BY EVALUATING THE POLYNOMIAL THERE. INFORMATION DEFINING THIS 9100
C POLYNOMIAL IS PASSED FROM STEP SO INTRP CANNOT BE USED ALONE. 9110
C 9120
C THIS CODE IS COMPLETELY EXPLAINED AND DOCUMENTED IN THE TEXT, 9130
C COMPUTER SOLUTION OF ORDINARY DIFFERENTIAL EQUATIONS& THE INITIAL 9140
C VALUE PROBLEM BY L. F. SHAMPINE AND M. K. GORDON. 9150
C 9160
C INPUT TO INTRP -- 9170
C 9180
C THE USER PROVIDES STORAGE IN THE CALLING PROGRAM FOR THE ARRAYS IN 9190
C THE CALL LIST 9200
C DIMENSION Y(NEQN),YOUT(NEQN),YPOUT(NEQN),PHI(NEQN,16),PSI(12) 9210
C AND DEFINES 9220
C XOUT -- POINT AT WHICH SOLUTION IS DESIRED. 9230
C THE REMAINING PARAMETERS ARE DEFINED IN STEP AND PASSED TO INTRP 9240
C FROM THAT SUBROUTINE 9250
C 9260
C OUTPUT FROM INTRP -- 9270
C 9280
C YOUT(*) -- SOLUTION AT XOUT 9290
C YPOUT(*) -- DERIVATIVE OF SOLUTION AT XOUT 9300
C THE REMAINING PARAMETERS ARE RETURNED UNALTERED FROM THEIR INPUT 9310
C VALUES. INTEGRATION WITH STEP MAY BE CONTINUED. 9320
C 9330
C DIMENSION G(13),W(13),RHO(13) 9340
C DATA G(1)/1.0/,RHO(1)/1.0/ 9350
C 9360
C HI = XOUT - X 9370
C KI = KOLD + 1 9380
C KIP1 = KI + 1 9390
C 9400

```

```

C   INITIALIZE W(*) FOR COMPUTING G(*)          9410
C   9420
C       DO 5 I = 1,KI                          9430
C           TEMP1 = I                          9440
C       5   W(I) = 1.0/TEMP1                    9450
C           TERM = 0.0                          9460
C   9470
C   COMPUTE G(*)                               9480
C   9490
C       DO 15 J = 2,KI                          9500
C           JM1 = J - 1                        9510
C           PSIJM1 = PSI(JM1)                 9520
C           GAMMA = (HI + TERM)/PSIJM1        9530
C           ETA = HI/PSIJM1                   9540
C           LIMIT1 = KIP1 - J                 9550
C           DO 10 I = 1,LIMIT1                9560
C       10  W(I) = GAMMA*W(I) - ETA*W(I+1)    9570
C           G(J) = W(1)                       9580
C           RHO(J) = GAMMA*RHO(JM1)           9590
C       15  TERM = PSIJM1                     9600
C   9610
C   INTERPOLATE                                9620
C   9630
C       DO 20 L = 1,NEQN                       9640
C           YPOUT(L) = 0.0                    9650
C       20  YOUT(L) = 0.0                      9660
C           DO 30 J = 1,KI                     9670
C               I = KIP1 - J                   9680
C               TEMP2 = G(I)                   9690
C               TEMP3 = RHO(I)                 9700
C           DO 25 L = 1,NEQN                   9710
C               YOUT(L) = YOUT(L) + TEMP2*PHI(L,I) 9720
C       25  YPOUT(L) = YPOUT(L) + TEMP3*PHI(L,I) 9730
C       30  CONTINUE                          9740
C           DO 35 L = 1,NEQN                   9750
C       35  YOUT(L) = Y(L) + HI*YOUT(L)        9760
C           RETURN                             9770
C       END                                    9780

SUBROUTINE ROOT(T,FT,B,C,RELERR,ABSERR,IFLAG) 9790
C   9800
C   ROOT COMPUTES A ROOT OF THE NONLINEAR EQUATION F(X)=0 9810
C   WHERE F(X) IS A CONTINUOUS REAL FUNCTION OF A SINGLE REAL 9820
C   VARIABLE X. THE METHOD USED IS A COMBINATION OF BISECTION 9830
C   AND THE SECANT RULE. 9840
C   9850
C   NORMAL INPUT CONSISTS OF A CONTINUOUS FUNCTION F AND AN 9860
C   INTERVAL (B,C) SUCH THAT F(B)*F(C).LE.0.0. EACH ITERATION 9870
C   FINDS NEW VALUES OF B AND C SUCH THAT THE INTERVAL(B,C) IS 9880
C   SHRUNK AND F(B)*F(C).LE.0.0. THE STOPPING CRITERION IS 9890
C   9900
C       ABS(B-C).LE.2.0*(RELERR*ABS(B)+ABSERR) 9910
C   9920
C   WHERE RELERR=RELATIVE ERROR AND ABSERR=ABSOLUTE ERROR ARE 9930
C   INPUT QUANTITIES. SET THE FLAG, IFLAG, POSITIVE TO INITIALIZE 9940
C   THE COMPUTATION. AS B,C AND IFLAG ARE USED FOR BOTH INPUT AND 9950
C   OUTPUT, THEY MUST BE VARIABLES IN THE CALLING PROGRAM. 9960
C   9970
C   IF 0 IS A POSSIBLE ROOT, ONE SHOULD NOT CHOOSE ABSERR=0.0. 9980
C   9990
C   THE OUTPUT VALUE OF B IS THE BETTER APPROXIMATION TO A ROOT 10000
C   AS B AND C ARE ALWAYS REDEFINED SO THAT ABS(F(B)).LE.ABS(F(C)). 10010
C   10020
C   TO SOLVE THE EQUATION, ROOT MUST EVALUATE F(X) REPEATEDLY. THIS 10030
C   IS DONE IN THE CALLING PROGRAM. WHEN AN EVALUATION OF F IS 10040
C   NEEDED AT T, ROOT RETURNS WITH IFLAG NEGATIVE. EVALUATE FT=F(T) 10050
C   AND CALL ROOT AGAIN. DO NOT ALTER IFLAG. 10060
C   10070
C   WHEN THE COMPUTATION IS COMPLETE, ROOT RETURNS TO THE CALLING 10080
C   PROGRAM WITH IFLAG POSITIVE= 10090
C   10100
C       IFLAG=1 IF F(B)*F(C).LT.0 AND THE STOPPING CRITERION IS MET. 10110
C   10120
C       =2 IF A VALUE B IS FOUND SUCH THAT THE COMPUTED VALUE 10130
C       F(B) IS EXACTLY ZERO. THE INTERVAL (B,C) MAY NOT 10140

```



```

P=(B-A)*FB 10820
Q=FA-FB 10830
IF(P.GE.0.0)GO TO 3 10840
P=-P 10850
Q=-Q 10860
C 10870
C UPDATE A, CHECK IF REDUCTION IN THE SIZE OF BRACKETING 10880
C INTERVAL IS SATISFACTORY. IF NOT BISECT UNTIL IT IS. 10890
C 10900
3 A=B 10910
FA=FB 10920
IC=IC+1 10930
IF(IC.LT.4)GO TO 4 10940
IF(8.0*ACMB.GE.ACBS)GO TO 6 10950
IC=0 10960
ACBS=ACMB 10970
C 10980
C TEST FOR TOO SMALL A CHANGE. 10990
C 11000
4 IF(P.GT.ABS(Q)*TOL)GO TO 5 11010
C 11020
C INCREMENT BY TOLERANCE 11030
C 11040
B=B+SIGN(TOL,CMB) 11050
GO TO 7 11060
C 11070
C ROOT OUGHT TO BE BETWEEN B AND (C+B)/2 11080
C 11090
5 IF(P.GE.CMB*Q)GO TO 6 11100
C 11110
C USE SECANT RULE. 11120
C 11130
B=B+P/Q 11140
GO TO 7 11150
C 11160
C USE BISECTION. 11170
C 11180
6 B=0.5*(C+B) 11190
C 11200
C HAVE COMPLETED COMPUTATION FOR NEW ITERATE B. 11210
C 11220
7 T=B 11230
IFLAG=-3 11240
RETURN 11250
400 FB=FT 11260
IF(FB.EQ.0.0)GO TO 9 11270
KOUNT=KOUNT+1 11280
IF(SIGN(1.0,FB).NE.SIGN(1.0,FC))GO TO 1 11290
C=A 11300
FC=FA 11310
GO TO 1 11320
C 11330
C FINISHED. SET IFLAG. 11340
C 11350
8 IF(SIGN(1.0,FB).EQ.SIGN(1.0,FC))GO TO 11 11360
IF(ABS(FB).GT.FX)GO TO 10 11370
IFLAG=1 11380
RETURN 11390
9 IFLAG=2 11400
RETURN 11410
10 IFLAG=3 11420
RETURN 11430
11 IFLAG=4 11440
RETURN 11450
12 IFLAG=5 11460
RETURN 11470
END 11480

```


ALGORITHM 556

Exponential Integrals [S13]

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Key Words and Phrases: exponential integrals, Miller algorithm, confluent hypergeometric functions

CR Categories: 5.12

Language: Fortran

DESCRIPTION

The Fortran subroutine EXPINT given here is an implementation of [1]. EXPINT has four machine-dependent parameters XCUT, XLIM, ETOL, and EULER which are set into DATA statements. XCUT is a breakpoint such that for $x \leq \text{XCUT}$, the series is evaluated and for $x > \text{XCUT}$, the Miller algorithm is applied. XLIM is the approximate underflow limit for e^{-x} , $x \geq 0$, and ETOL is nominally set to $1.E-D$ where D is the number of base 10 digits in a word. EULER is the negative Euler constant, $-\gamma$. Thus

DATA XCUT, XLIM, ETOL/2.0, 667.0, 1.E-12/

would be appropriate for CDC single-precision arithmetic, while

DATA XCUT, XLIM, ETOL/1.0, 172.0, 1.E-6/

would be appropriate for IBM single-precision arithmetic. The two choices for XCUT reflect the fact that there is a loss of up to two digits on $1 < x \leq 2$ with the series evaluation. This loss can be tolerated on longer word length machines, but not on shorter word length machines.

Maximum accuracy can always be achieved with $\text{XCUT} = 1$. However, for longer word lengths where $D = 14$, the reduction in computation by moving XCUT from 1 to 2, achieved at the expense of two digits of accuracy, seems to be a worthwhile trade-off, and $D = 12$ reflects this modification for CDC machines. $D = 12$ is also more consistent with the accuracy attainable from $\text{EXP}(-X)$ near the underflow limit $X = 667$.

While the subroutine EXPINT is almost portable, the function DIGAM, which computes the psi function at integer arguments, is supplied as a CDC 6600-7600 Fortran function. Modifications for other machines can be made easily from eq. (7) of [1].

The convergence of eq. (2) in [1] is so rapid that the $m = n - 1$ term, which requires $\psi(n)$, is reached for only small values of n . A table of 100 values suffices for virtually all single-precision implementations of EXPINT. An initialization step to generate a higher precision table might be appropriate if multiple precision is anticipated. For CDC single precision or IBM double precision, 36 values suffice for n as high as 10^{12} and relative errors of 10^{-14} .

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Testing of EXPINT

Extensive testing of EXPINT against a quadrature

$$E_n(x) = e^{-x} x^{n-1} \int_x^\infty \frac{e^{-(t-x)}}{t^n} dt$$

was used to check out the routine on parameter ranges

$$E_n(x) = e^{-x} x^{n-1} \int_x^\infty \frac{e^{-(t-x)}}{t^n} dt$$

was used to check out the routine on parameter ranges

$$\begin{aligned} 10^{-11} &\leq \text{TOL} \leq 0.1 \\ 0 &\leq X \leq 100 \\ 1 &\leq N \leq 100 \\ 1 &\leq M \leq 100 \\ \text{XCUT} &= 1, 2 \\ \text{KODE} &= 1, 2 \end{aligned}$$

where KODE is used to scale $E_{N+k}(x)$, $k = 0, 1, \dots, M-1$ by e^x , if desired. Some spot checks for x and N in the hundreds, as well as special values $\text{XCUT} \pm 10^{-13}$, were also made at tolerances TOL as low as 10^{-9} . The quadrature could be relied upon to give the requested accuracy down to $\text{TOL} = 10^{-11}$ over the full exponent range of the CDC 6600. For large x , argument reduction in computing e^{-x} will result in decreased accuracy, up to three digits near the underflow limit $x \sim 667$. Hence tolerances smaller than 10^{-11} would require double-precision exponentials to obtain checks for large x .

Relative errors in long sequences, up to 100 terms, were checked for degradation of the recurrence by comparing members of a sequence with single evaluations ($M = 1$) at high accuracies. When $M = 1$, EXPINT uses one of the two basic methods for evaluation rather than recursion.

Relative errors in all of these tests were less than the relative error specified in EXPINT to generate the test values.

REFERENCE

1. Amos, D.E. Computation of exponential integrals. *ACM Trans. Math. Softw.* 6, 3 (Sept. 1980), 365-377.

ALGORITHM

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SUBROUTINE EXPINT(X, N, KODE, M, TOL, EN, IERR)          EXP 10
C                                                       EXP 20
C WRITTEN BY D.E. AMOS, SANDIA LABORATORIES, ALBUQUERQUE, NM, 87185 EXP 30
C                                                       EXP 40
C REFERENCE                                             EXP 50
C   COMPUTATION OF EXPONENTIAL INTEGRALS BY D.E. AMOS, ACM EXP 60
C   TRANS. MATH SOFTWARE, 1980                          EXP 70
C                                                       EXP 80
C ABSTRACT                                             EXP 90
C   EXPINT COMPUTES M MEMBER SEQUENCES OF EXPONENTIAL INTEGRALS EXP 100
C   E(N+K,X), K=0,1,...,M-1 FOR N.GE.1 AND X.GE.0. THE POWER EXP 110
C   SERIES IS IMPLEMENTED FOR X.LE.XCUT AND THE CONFLUENT EXP 120
C   HYPERGEOMETRIC REPRESENTATION                     EXP 130
C                                                       EXP 140
C   E(A,X) = EXP(-X)*(X**(A-1))*U(A,A,X)                EXP 150
C                                                       EXP 160
C   IS COMPUTED FOR X.GT.XCUT. SINCE SEQUENCES ARE COMPUTED IN A EXP 170
C   STABLE FASHION BY RECURRING AWAY FROM X, A IS SELECTED AS THE EXP 180
C   INTEGER CLOSEST TO X WITHIN THE CONSTRAINT N.LE.A.LE.N+M-1. EXP 190
C   FOR THE U COMPUTATION A IS FURTHER MODIFIED TO BE THE EXP 200
C   NEAREST EVEN INTEGER. INDICES ARE CARRIED FORWARD OR EXP 210
C   BACKWARD BY THE TWO TERM RECURSION RELATION       EXP 220
C                                                       EXP 230
C   K*E(K+1,X) + X*E(K,X) = EXP(-X)                   EXP 240

```

```

C
C      ONCE E(A,X) IS COMPUTED. THE U FUNCTION IS COMPUTED BY MEANS EXP 250
C      OF THE BACKWARD RECURSIVE MILLER ALGORITHM APPLIED TO THE EXP 260
C      THREE TERM CONTIGUOUS RELATION FOR U(A+K,A,X), K=0,1,... EXP 270
C      THIS PRODUCES ACCURATE RATIOS AND DETERMINES U(A+K,A,X),AND EXP 280
C      HENCE E(A,X), TO WITHIN A MULTIPLICATIVE CONSTANT C. EXP 290
C      ANOTHER CONTIGUOUS RELATION APPLIED TO C*U(A,A,X) AND EXP 300
C      C*U(A+1,A,X) GETS C*U(A+1,A+1,X), A QUANTITY PROPORTIONAL TO EXP 310
C      E(A+1,X). THE NORMALIZING CONSTANT C IS OBTAINED FROM THE EXP 320
C      TWO TERM RECURSION RELATION ABOVE WITH K=A. EXP 330
C      EXP 340
C      EXP 350
C      MACHINE DEPENDENT PARAMETERS - XCUT, XLIM, ETOL, EULER, DIGAM EXP 360
C      EXP 370
C      EXPINT WRITES ERROR DIAGNOSTICS TO LOGICAL UNIT 3 EXP 380
C      EXP 390
C      DESCRIPTION OF ARGUMENTS EXP 400
C      EXP 410
C      INPUT EXP 420
C      X      X.GT.0.0 FOR N=1 AND X.GE.0.0 FOR N.GE.2 EXP 430
C      N      ORDER OF THE FIRST MEMBER OF THE SEQUENCE, N.GE.1 EXP 440
C      KODE   A SELECTION PARAMETER FOR SCALED VALUES EXP 450
C      KODE=1 RETURNS E(N+K,X), K=0,1,...,M-1. EXP 460
C      =2 RETURNS EXP(X)*E(N+K,X), K=0,1,...,M-1. EXP 470
C      M      NUMBER OF EXPONENTIAL INTEGRALS IN THE SEQUENCE, EXP 480
C      M.GE.1 EXP 490
C      TOL    RELATIVE ACCURACY WANTED, ETOL.LE.TOL.LE.0.1 EXP 500
C      ETOL=1.E-12 EXP 510
C      EXP 520
C      OUTPUT EXP 530
C      EN     A VECTOR OF DIMENSION AT LEAST M CONTAINING VALUES EXP 540
C      EN(K) = E(N+K-1,X) OR EXP(X)*E(N+K-1,X), K=1,M EXP 550
C      DEPENDING ON KODE EXP 560
C      IERR   UNDERFLOW INDICATOR EXP 570
C      IERR=0 A NORMAL RETURN EXP 580
C      =1 X EXCEEDS XLIM AND AN UNDERFLOW OCCURS. EXP 590
C      EN(K)=0.0, K=1,M RETURNED ON KODE=1 EXP 600
C      XLIM=667. EXP 610
C      EXP 620
C      ERROR CONDITIONS EXP 630
C      AN IMPROPER INPUT PARAMETER IS A FATAL ERROR EXP 640
C      UNDERFLOW IS A NON FATAL ERROR. ZERO ANSWERS ARE RETURNED. EXP 650
C      EXP 660
C      DIMENSION EN(1), A(99), B(99), Y(2) EXP 670
C      EXP 680
C      DATA XCUT, XLIM, ETOL /2.0E0,667.0E0,1.0E-12/ EXP 690
C      DATA EULER /-5.77215664901533E-01/ EXP 700
C      DATA LUN /3/ EXP 710
C      EXP 720
C      IF (N.LT.1) GO TO 260 EXP 730
C      IF (KODE.LT.1 .OR. KODE.GT.2) GO TO 270 EXP 740
C      IF (M.LT.1) GO TO 280 EXP 750
C      IF (TOL.LT.ETOL .OR. TOL.GT.0.1E0) GO TO 290 EXP 760
C      EXP 770
C      IERR = 0 EXP 780
C      IF (X.GT.XCUT) GO TO 100 EXP 790
C      IF (X.LT.0.0E0) GO TO 300 EXP 800
C      IF (X.EQ.0.0E0 .AND. N.EQ.1) GO TO 310 EXP 810
C      IF (X.EQ.0.0E0 .AND. N.GT.1) GO TO 80 EXP 820
C      EXP 830
C      SERIES FOR E(N,X) FOR X.LE.XCUT EXP 840
C      EXP 850
C      IX = INT(X+0.5E0) EXP 860
C      ICASE=1 MEANS INTEGER CLOSEST TO X IS 2 AND N=1 EXP 870
C      ICASE=2 MEANS INTEGER CLOSEST TO X IS 0,1, OR 2 AND N.GE.2 EXP 880
C      ICASE = 2 EXP 890
C      IF (IX.GT.N) ICASE = 1 EXP 900
C      NM = N - ICASE + 1 EXP 910
C      ND = NM + 1 EXP 920
C      IND = 3 - ICASE EXP 930
C      MU = M - IND EXP 940
C      ML = 1 EXP 950
C      KS = ND EXP 960
C      FNM = FLOAT(NM) EXP 970
C      S = 0.0E0 EXP 980
C      XTOL = 3.0E0*TOL EXP 990
C      IF (ND.EQ.1) GO TO 100 EXP 1000
C      XTOL = 0.3333E0*TOL EXP 1010

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```

      S = 1.0E0/FNM
10 CONTINUE
      AA = 1.0E0
      AK = 1.0E0
      DO 50 I=1,35
        AA = -AA*X/AK
        IF (I.EQ.NM) GO TO 30
        S = S - AA/(AK-FNM)
        IF (ABS(AA).LE.XTOL*ABS(S)) GO TO 20
        AK = AK + 1.0E0
        GO TO 50
20 CONTINUE
      IF (I.LT.2) GO TO 40
      IF (ND-2.GT.I .OR. I.GT.ND-1) GO TO 60
      AK = AK + 1.0E0
      GO TO 50
30 S = S + AA*(-ALOG(X)+DIGAM(ND))
      XTOL = 3.0E0*TOL
40 AK = AK + 1.0E0
50 CONTINUE
      GO TO 320
60 IF (ND.EQ.1) S = S + (-ALOG(X)+EULER)
      IF (KODE.EQ.2) S = S*EXP(X)
      EN(1) = S
      EMX = 1.0E0
      IF (M.EQ.1) GO TO 70
      EN(IND) = S
      AA = FLOAT(KS)
      IF (KODE.EQ.1) EMX = EXP(-X)
      GO TO (220, 240), ICASE
70 IF (ICASE.EQ.2) RETURN
      IF (KODE.EQ.1) EMX = EXP(-X)
      EN(1) = (EMX-S)/X
      RETURN
80 CONTINUE
      DO 90 I=1,M
        EN(I) = 1.0E0/FLOAT(N+I-2)
90 CONTINUE
      RETURN
C
C   BACKWARD RECURSIVE MILLER ALGORITHM FOR
C       E(N,X)=EXP(-X)*(X**(N-1))*U(N,N,X)
C   WITH RECURSION AWAY FROM N=INTEGER CLOSEST TO X.
C   U(A,B,X) IS THE SECOND CONFLUENT HYPERGEOMETRIC FUNCTION
C
100 CONTINUE
      EMX = 1.0E0
      IF (KODE.EQ.2) GO TO 130
      IF (X.LE.XLIM) GO TO 120
      IERR = 1
      DO 110 I=1,M
        EN(I) = 0.0E0
110 CONTINUE
      RETURN
120 EMX = EXP(-X)
130 CONTINUE
      IX = INT(X+0.5E0)
      KN = N + M - 1
      IF (KN.LE.IX) GO TO 140
      IF (N.LT.IX .AND. IX.LT.KN) GO TO 170
      IF (N.GE.IX) GO TO 160
      GO TO 340
140 ICASE = 1
      KS = KN
      ML = M - 1
      MU = -1
      IND = M
      IF (KN.GT.1) GO TO 180
150 KS = 2
      ICASE = 3
      GO TO 180
160 ICASE = 2
      IND = 1
      KS = N
      MU = M - 1
      IF (N.GT.1) GO TO 180

```

```

EXP 1020
EXP 1030
EXP 1040
EXP 1050
EXP 1060
EXP 1070
EXP 1080
EXP 1090
EXP 1100
EXP 1110
EXP 1120
EXP 1130
EXP 1140
EXP 1150
EXP 1160
EXP 1170
EXP 1180
EXP 1190
EXP 1200
EXP 1210
EXP 1220
EXP 1230
EXP 1240
EXP 1250
EXP 1260
EXP 1270
EXP 1280
EXP 1290
EXP 1300
EXP 1310
EXP 1320
EXP 1330
EXP 1340
EXP 1350
EXP 1360
EXP 1370
EXP 1380
EXP 1390
EXP 1400
EXP 1410
EXP 1420
EXP 1430
EXP 1440
EXP 1450
EXP 1460
EXP 1470
EXP 1480
EXP 1490
EXP 1500
EXP 1510
EXP 1520
EXP 1530
EXP 1540
EXP 1550
EXP 1560
EXP 1570
EXP 1580
EXP 1590
EXP 1600
EXP 1610
EXP 1620
EXP 1630
EXP 1640
EXP 1650
EXP 1660
EXP 1670
EXP 1680
EXP 1690
EXP 1700
EXP 1710
EXP 1720
EXP 1730
EXP 1740
EXP 1750
EXP 1760
EXP 1770

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```

      IF (KN.EQ.1) GO TO 150
      IX = 2
170 ICASE = 1
      KS = IX
      ML = IX - N
      IND = ML + 1
      MU = KN - IX
180 CONTINUE
      IK = KS/2
      AH = FLOAT(IK)
      JSET = 1 + KS - (IK+IK)
C     START COMPUTATION FOR
C           EN(IND) = C*U( A , A ,X)      JSET=1
C           EN(IND) = C*U(A+1,A+1,X)      JSET=2
C     FOR AN EVEN INTEGER A.
      IC = 0
      AA = AH + AH
      AAMS = AA - 1.0E0
      AAMS = AAMS*AAMS
      TX = X + X
      FX = TX + TX
      AK = AH
      XTOL = TOL
      IF (TOL.LE.1.0E-3) XTOL = 20.0E0*TOL
      CT = AAMS + FX*AH
      EM = (AH+1.0E0)/((X+AA)*XTOL*SQRT(CT))
      BK = AA
      CC = AH*AH
C     FORWARD RECURSION FOR P(IC),P(IC+1) AND INDEX IC FOR BACKWARD
C     RECURSION
      P1 = 0.0E0
      P2 = 1.0E0
190 CONTINUE
      IF (IC.EQ.99) GO TO 330
      IC = IC + 1
      AK = AK + 1.0E0
      AT = BK/(BK+AK+CC+FLOAT(IC))
      BK = BK + AK + AK
      A(IC) = AT
      BT = (AK+AK+X)/(AK+1.0E0)
      B(IC) = BT
      PT = P2
      P2 = BT*P2 - AT*P1
      P1 = PT
      CT = CT + FX
      EM = EM*AT*(1.0E0-TX/CT)
      IF (EM*(AK+1.0E0).GT.P1*P1) GO TO 190
      ICT = IC
      KK = IC + 1
      BT = TX/(CT+FX)
      Y2 = (BK/(BK+CC+FLOAT(KK)))*(P1/P2)*(1.0E0-BT+0.375E0*BT*BT)
      Y1 = 1.0E0
C     BACKWARD RECURRENCE FOR
C           Y1= C*U( A ,A,X)
C           Y2= C*(A/(1+A/2))*U(A+1,A,X)
C     DO 200 K=1,ICT
      KK = KK - 1
      YT = Y1
      Y1 = (B(KK)*Y1-Y2)/A(KK)
      Y2 = YT
200 CONTINUE
C     THE CONTIGUOUS RELATION
C           X*U(B,C+1,X)=(C-B)*U(B,C,X)+U(B-1,C,X)
C     WITH B=A+1 , C=A IS USED FOR
C           Y(2) = C * U(A+1,A+1,X)
C     X IS INCORPORATED INTO THE NORMALIZING RELATION FOR CNORM.
      Y(1) = Y1
      Y(2) = Y1 - Y2*(AH+1.0E0)/AA
      CNORM = EMX/(AA*Y(2)+X*Y(1))
      IF (ICASE.EQ.3) GO TO 210
      EN(IND) = CNORM*Y(JSET)
      IF (M.EQ.1) RETURN
      AA = FLOAT(KS)
      GO TO (220, 240), ICASE
C
C     RECURSION SECTION N*E(N+1,X) + X*E(N,X)=EMX
C

```

```

EXP 1780
EXP 1790
EXP 1800
EXP 1810
EXP 1820
EXP 1830
EXP 1840
EXP 1850
EXP 1860
EXP 1870
EXP 1880
EXP 1890
EXP 1900
EXP 1910
EXP 1920
EXP 1930
EXP 1940
EXP 1950
EXP 1960
EXP 1970
EXP 1980
EXP 1990
EXP 2000
EXP 2010
EXP 2020
EXP 2030
EXP 2040
EXP 2050
EXP 2060
EXP 2070
EXP 2080
EXP 2090
EXP 2100
EXP 2110
EXP 2120
EXP 2130
EXP 2140
EXP 2150
EXP 2160
EXP 2170
EXP 2180
EXP 2190
EXP 2200
EXP 2210
EXP 2220
EXP 2230
EXP 2240
EXP 2250
EXP 2260
EXP 2270
EXP 2280
EXP 2290
EXP 2300
EXP 2310
EXP 2320
EXP 2330
EXP 2340
EXP 2350
EXP 2360
EXP 2370
EXP 2380
EXP 2390
EXP 2400
EXP 2410
EXP 2420
EXP 2430
EXP 2440
EXP 2450
EXP 2460
EXP 2470
EXP 2480
EXP 2490
EXP 2500
EXP 2510
EXP 2520
EXP 2530
EXP 2540

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```

210 EN(1) = (EMX-CNORM*Y(1))/X      EXP 2550
    RETURN                          EXP 2560
220 K = IND - 1                      EXP 2570
    DO 230 I=1,ML                    EXP 2580
        AA = AA - 1.0E0              EXP 2590
        EN(K) = (EMX-AA*EN(K+1))/X   EXP 2600
        K = K - 1                    EXP 2610
230 CONTINUE                         EXP 2620
    IF (MU.LE.0) RETURN              EXP 2630
    AA = FLOAT(KS)                   EXP 264
240 K = IND                          EXP 2650
    DO 250 I=1,MU                    EXP 2660
        EN(K+1) = (EMX-X*EN(K))/AA  EXP 2670
        AA = AA + 1.0E0              EXP 2680
        K = K + 1                    EXP 2690
250 CONTINUE                         EXP 2700
    RETURN                            EXP 2710
C                                    EXP 2720
C                                    EXP 2730
260 WRITE (LUN,99999)                EXP 2740
    RETURN                            EXP 2750
270 WRITE (LUN,99998)                EXP 2760
    RETURN                            EXP 2770
280 WRITE (LUN,99997)                EXP 2780
    RETURN                            EXP 2790
290 WRITE (LUN,99996)                EXP 2800
    RETURN                            EXP 2810
300 WRITE (LUN,99995)                EXP 2820
    RETURN                            EXP 2830
310 WRITE (LUN,99994)                EXP 2840
    RETURN                            EXP 2850
320 WRITE (LUN,99993)                EXP 2860
    RETURN                            EXP 2870
330 WRITE (LUN,99992)                EXP 2880
    RETURN                            EXP 2890
340 WRITE (LUN,99991)                EXP 2900
    RETURN                            EXP 2910
99999 FORMAT (32H IN EXPINT, N NOT GREATER THAN 0) EXP 2920
99998 FORMAT (27H IN EXPINT, KODE NOT 1 OR 2)     EXP 2930
99997 FORMAT (32H IN EXPINT, M NOT GREATER THAN 0) EXP 2940
99996 FORMAT (33H IN EXPINT, TOL NOT WITHIN LIMITS) EXP 2950
99995 FORMAT (37H IN EXPINT, X IS NOT ZERO OR POSITIVE) EXP 2960
99994 FORMAT (46H IN EXPINT, THE EXPONENTIAL INTEGRAL IS NOT DE, EXP 2970
    * 21HFINED FOR X=0 AND N=1)                   EXP 2980
99993 FORMAT (46H IN EXPINT, RELATIVE ERROR TEST FOR SERIES TER, EXP 2990
    * 28HMINATION NOT MET IN 36 TERMS)           EXP 3000
99992 FORMAT (46H IN EXPINT, TERMINATION TEST FOR MILLER ALGORI, EXP 3010
    * 23HTHM NOT MET IN 99 STEPS)               EXP 3020
99991 FORMAT (46H IN EXPINT, AN ERROR IN PLACING INT(X+0.5) WIT, EXP 3030
    * 47HH RESPECT TO N AND N+M-1 OCCURRED FOR X.GT.XCUT) EXP 3040
    END                                        EXP 3050

    FUNCTION DIGAM(N)                      DIG 10
C                                    DIG 20
C    THIS SUBROUTINE RETURNS VALUES OF PSI(X)=DERIVATIVE OF LOG DIG 30
C    GAMMA(X), X.GT.0.0 AT INTEGER ARGUMENTS. A TABLE LOOK-UP IS DIG 40
C    PERFORMED FOR N.LE.100, AND THE ASYMPTOTIC EXPANSION IS DIG 50
C    EVALUATED FOR N.GT.100.              DIG 60
C                                    DIG 70
C    DIMENSION B(4), C(100), C1(32), C2(27), C3(22), C4(19) DIG 80
C    EQUIVALENCE (C(1),C1(1))              DIG 90
C    EQUIVALENCE (C(33),C2(1))            DIG 100
C    EQUIVALENCE (C(60),C3(1))            DIG 110
C    EQUIVALENCE (C(82),C4(1))            DIG 120
C                                    DIG 130
C    DATA C1 /-5.7721566490153E-01,4.22784335098467E-01, DIG 140
C    * 9.22784335098467E-01,1.25611766843180E+00,1.50611766843180E+00, DIG 150
C    * 1.70611766843180E+00,1.87278433509847E+00,2.01564147795561E+00, DIG 160
C    * 2.14064147795561E+00,2.25175258906672E+00,2.35175258906672E+00, DIG 170
C    * 2.44266167997581E+00,2.52599501330915E+00,2.60291809023222E+00, DIG 180
C    * 2.67434666166079E+00,2.74101332832746E+00,2.80351332832746E+00, DIG 190
C    * 2.86233685773923E+00,2.91789241329478E+00,2.97052399224215E+00, DIG 200
C    * 3.02052399224215E+00,3.06814303986120E+00,3.11359758531574E+00, DIG 210
C    * 3.15707584618531E+00,3.19874251285197E+00,3.23874251285197E+00, DIG 220
C    * 3.27720005131351E+00,3.31424108835055E+00,3.34995537406484E+00, DIG 230

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* 3.38443813268552E+00, 3.41777146601886E+00, 3.45002953053499E+00/ DIG 240
DATA C2 /3.48127953053499E+00, 3.51158256083802E+00, DIG 250
* 3.54099432554390E+00, 3.56956575411533E+00, 3.59734353189311E+00, DIG 260
* 3.62437055892013E+00, 3.65068634839382E+00, 3.67632737403484E+00, DIG 270
* 3.70132737403484E+00, 3.72571761793728E+00, 3.74952714174681E+00, DIG 280
* 3.77278295570029E+00, 3.79551022842757E+00, 3.81773245064979E+00, DIG 290
* 3.83947158108457E+00, 3.86074817682925E+00, 3.88158151016259E+00, DIG 300
* 3.90198967342789E+00, 3.92198967342789E+00, 3.94159751656515E+00, DIG 310
* 3.96082828579592E+00, 3.97969621032422E+00, 3.99821472884274E+00, DIG 320
* 4.01639654702455E+00, 4.03425368988170E+00, 4.05179754953082E+00, DIG 330
* 4.06903892884117E+00/ DIG 340
DATA C3 /4.08598808138354E+00, 4.10265474805020E+00, DIG 350
* 4.11904819067316E+00, 4.13517722293122E+00, 4.15105023880424E+00, DIG 360
* 4.16667523880424E+00, 4.18205985418885E+00, 4.19721136934037E+00, DIG 370
* 4.21213674247470E+00, 4.22684262482764E+00, 4.24133537845082E+00, DIG 380
* 4.25562109273654E+00, 4.26970559977879E+00, 4.28359448866768E+00, DIG 390
* 4.29729311880467E+00, 4.31080663231818E+00, 4.32413996565151E+00, DIG 400
* 4.33729786038836E+00, 4.35028487337537E+00, 4.36310538619588E+00, DIG 410
* 4.37576361404398E+00, 4.38826361404398E+00/ DIG 420
DATA C4 /4.40060929305633E+00, 4.41280441500755E+00, DIG 430
* 4.42485260777863E+00, 4.43675736968340E+00, 4.44852207556575E+00, DIG 440
* 4.46014998254249E+00, 4.47164423541606E+00, 4.48300787177969E+00, DIG 450
* 4.49424382683587E+00, 4.50535493794698E+00, 4.51634394893599E+00, DIG 460
* 4.52721351415338E+00, 4.53796620232543E+00, 4.54860450019777E+00, DIG 470
* 4.55913081598724E+00, 4.56954748265391E+00, 4.57985676100442E+00, DIG 480
* 4.59006084263708E+00, 4.60016185273809E+00/ DIG 490
C DIG 500
DATA B /1.66666666666667E-01, -3.33333333333333E-02, DIG 510
* 2.38095238095238E-02, -3.33333333333333E-02/ DIG 520
C DIG 530
IF (N.GT.100) GO TO 10 DIG 540
DIGAM = C(N) DIG 550
RETURN DIG 560
10 FN = N DIG 570
AX = 1.0E0 DIG 580
AK = 2.0E0 DIG 590
S = -0.5E0/FN DIG 600
IF (FN.GT.1.E+8) GO TO 30 DIG 610
FN2 = FN*FN DIG 620
DO 20 K=1, 3 DIG 630
AX = AX*FN2 DIG 640
S = S - B(K)/(AX*AK) DIG 650
AK = AK + 2.0E0 DIG 660
20 CONTINUE DIG 670
30 CONTINUE DIG 680
DIGAM = S + ALOG(FN) DIG 690
RETURN DIG 700
END DIG 710

C PROGRAM TSTEXP(INPUT,OUTPUT,TAPE3=OUTPUT) 00000010
C 00000020
C PROGRAM TO TEST SUBROUTINE EXPINT AGAINST AN ADAPTIVE QUADRATURE. 00000030
C PARAMETER VALUES ARE PRINTED AND, IN THE EVENT THAT THE RELATIVE 00000040
C ERROR TEST IS NOT SATISFIED, X, ERROR, N, AND KODE ARE ALSO 00000050
C PRINTED. AN OUTPUT WITH ONLY PARAMETER VALUES INDICATES THAT ALL 00000060
C TESTS WERE PASSED. GAUS8 COMPUTES THE QUADRATURES. 00000070
C 00000080
C DIMENSION XTOL(10), EN(50), EV(50) 00000090
IOUT = 3 00000100
NL = 1 00000110
NU = 16 00000120
NINC = 5 00000130
ML = 1 00000140
MU = 25 00000150
MINC = 8 00000160
KM = 5 00000170
JL = 1 00000180
JU = 40 00000190
JINC = 3 00000200
XTOL(1) = 1.0E-2 00000210
DO 10 I=2, 3 00000220
XTOL(I) = XTOL(I-1)*1.0E-3 00000230
10 CONTINUE 00000240
DO 80 IT=1, 3 00000250
TOL = XTOL(IT) 00000260

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TOLA = AMAX1(1.0E-12,TOL/10.0E0)      00000270
BTOL = TOL                             00000280
WRITE (IOUT,99999) TOL                 00000290
DO 70 M=ML,MU,MINC                     00000300
  WRITE (IOUT,99998) M                 00000310
  DO 60 N=NL,NU,NINC                   00000320
    WRITE (IOUT,99997) N               00000330
    DO 50 J=JL,JU,JINC                 00000340
      X = FLOAT(J-1)/5.0E0            00000350
      EX = EXP(-X)                    00000360
      IF (X.EQ.0. .AND. N.EQ.1) GO TO 50 00000370
      CALL EXPINT(X, N, 1, M, TOL, EN, IERR) 00000380
      CALL EXPINT(X, N, 2, M, TOL, EV, IERR) 00000390
      DO 40 K=1,M,KM                   00000400
        IF (X.GT.0.) GO TO 20         00000410
        IF (N+K.EQ.2) GO TO 40        00000420
        Y = 1.0E0/FLOAT(N+K-2)        00000430
        YY = Y                         00000440
        GO TO 30                       00000450
20      CONTINUE                       00000460
        NN = N + K - 1                 00000470
        YY = EINT(NN,X,TOLA,2)         00000480
        Y = YY*EX                      00000490
30      CONTINUE                       00000500
        ER = ABS((Y-EN(K))/Y)          00000510
        KODE = 1                       00000520
        IF (ER.GT.BTOL) WRITE (IOUT,99996) X, ER, NN, KODE 00000530
        KODE = 2                       00000540
        ERR = ABS((YY-EV(K))/YY)       00000550
        IF (ERR.GT.BTOL) WRITE (IOUT,99996) X, ERR, NN, KODE 00000560
40      CONTINUE                       00000570
50      CONTINUE                       00000580
60      CONTINUE                       00000590
70      CONTINUE                       00000600
80      CONTINUE                       00000610
      STOP                             00000620
99999 FORMAT (1H0, 5H TOL=, E15.4/)    00000630
99998 FORMAT (1H0, 2HM=, I5/)          00000640
99997 FORMAT (3X, 2HN=, I5)           00000650
99996 FORMAT (2E15.6, 2I5)            00000660
      END                               00000670

FUNCTION EINT(N, X, TOL, KODE)          00000680
COMMON /GEINT/ XX, FN                  00000690
EXTERNAL FEINT                         00000700
XX = X                                 00000710
FN = N                                 00000720
SIG = 1.0E0                            00000730
S = 0.0E0                               00000740
TOLA = TOL                              00000750
B = X                                   00000760
10 CONTINUE                             00000770
A = B                                   00000780
REL = TOL                               00000790
B = B + SIG                             00000800
CALL GAUS8(FEINT, A, B, REL, ANS, IERR) 00000810
S = S + ANS                             00000820
IF (ABS(ANS).LT.S*TOLA) GO TO 20        00000830
GO TO 10                                 00000840
20 EINT = S*EXP((FN-1.0E0)*ALOG(X)-FLOAT(2-KODE)*X) 00000850
RETURN                                  00000860
END                                      00000870

FUNCTION FEINT(T)                      00000880
COMMON /GEINT/ XX, FN                  00000890
FEINT = EXP(-T+XX-FN*ALOG(T))          00000900
RETURN                                  00000910
END                                      00000920

SUBROUTINE GAUS8 (FUN,A,B,ERR,ANS,IERR) 00000930
C                                       00000940
C BY RONDALL E JONES, SANDIA LABORATORIES 00000950
C SALIENT FEATURES -- INTERVAL BISECTION, COMBINED RELATIVE/ABSOLUTE 00000960

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C      ERROR CONTROL, COMPUTED MAXIMUM REFINEMENT LEVEL WHEN A IS      00000970
C      CLOSE TO B.                                                    00000980
C      00000990
C      ABSTRACT                                                         00001000
C      GAUS8 INTEGRATES REAL FUNCTIONS OF ONE VARIABLE OVER FINITE    00001010
C      INTERVALS, USING AN ADAPTIVE 8-POINT LEGENDRE-GAUSS ALGORITHM.  00001020
C      GAUS8 IS INTENDED PRIMARILY FOR HIGH ACCURACY INTEGRATION      00001030
C      OR INTEGRATION OF SMOOTH FUNCTIONS. FOR LOWER ACCURACY        00001040
C      INTEGRATION OF FUNCTIONS WHICH ARE NOT VERY SMOOTH,           00001050
C      EITHER QNC3 OR QNC7 MAY BE MORE EFFICIENT.                    00001060
C      00001070
C      DESCRIPTION OF ARGUMENTS                                         00001080
C      00001090
C      INPUT--                                                           00001100
C      FUN - NAME OF EXTERNAL FUNCTION TO BE INTEGRATED. THIS NAME    00001110
C      MUST BE IN AN EXTERNAL STATEMENT IN THE CALLING PROGRAM.      00001120
C      FUN MUST BE A FUNCTION OF ONE REAL ARGUMENT. THE VALUE        00001130
C      OF THE ARGUMENT TO FUN IS THE VARIABLE OF INTEGRATION        00001140
C      WHICH RANGES FROM A TO B.                                      00001150
C      A - LOWER LIMIT OF INTEGRAL                                     00001160
C      B - UPPER LIMIT OF INTEGRAL (MAY BE LESS THAN A)              00001170
C      ERR - IS A REQUESTED ERROR TOLERANCE. NORMALLY PICK A VALUE    00001180
C      ABS(ERR).LT.1.E-3. ANS WILL NORMALLY HAVE NO MORE ERROR      00001190
C      THAN ABS(ERR) TIMES THE INTEGRAL OF THE ABSOLUTE VALUE        00001200
C      OF FUN(X). USUALLY, SMALLER VALUES FOR ERR YIELD            00001210
C      MORE ACCURACY AND REQUIRE MORE FUNCTION EVALUATIONS.         00001220
C      A NEGATIVE VALUE FOR ERR CAUSES AN ESTIMATE OF THE           00001230
C      ABSOLUTE ERROR IN ANS TO BE RETURNED IN ERR.                  00001240
C      00001250
C      OUTPUT--                                                         00001260
C      ERR - WILL BE AN ESTIMATE OF THE ERROR IN ANS IF THE INPUT     00001270
C      VALUE OF ERR WAS NEGATIVE. THE ESTIMATED ERROR IS SOLELY      00001280
C      FOR INFORMATION TO THE USER AND SHOULD NOT BE USED AS        00001290
C      A CORRECTION TO THE COMPUTED INTEGRAL.                         00001300
C      ANS - COMPUTED VALUE OF INTEGRAL                                00001310
C      IERR- A STATUS CODE                                             00001320
C      --NORMAL CODES                                                 00001330
C      1 ANS MOST LIKELY MEETS REQUESTED ERROR TOLERANCE,           00001340
C      OR A=B.                                                         00001350
C      -1 A AND B ARE TOO NEARLY EQUAL TO ALLOW NORMAL              00001360
C      INTEGRATION. ANS IS SET TO ZERO.                               00001370
C      --ABNORMAL CODE                                               00001380
C      2 ANS PROBABLY DOES NOT MEET REQUESTED ERROR TOLERANCE.      00001390
C      00001400
C      00001410
C      00001420
C      GAUS8 USES SUBROUTINES ERRCHK, ERRGET, ERRPRT, ERXSET, ERSTGT  00001430
C      COMPILER DECKS GAUS8, ERRCHK                                   00001440
C      00001450
C      DIMENSION AA(30),HH(30),LR(30),VL(30),GR(30)                  00001460
C      DATA X1,X2,X3,X4/0.18343 46424 95650 , 0.52553 24099 16329 ,  00001470
C      1 0.79666 64774 13627 , 0.96028 98564 97536 /                00001480
C      DATA W1,W2,W3,W4/0.36268 37833 78362 , 0.31370 66458 77887 ,  00001490
C      1 0.22238 10344 53374 , 0.10122 85362 90376 /                00001500
C      DATA SQ2/1.41421356/,ICALL/0/                                00001510
C      DATA NLMN/1/,NLMX/30/,KMX/5000/,KML/6/,NBITS/48/           00001520
C      G8(X,H) = H*( (W1*(FUN(X-X1*H)+FUN(X+X1*H))                    00001530
C      1 +W2*(FUN(X-X2*H)+FUN(X+X2*H)))                               00001540
C      2 +(W3*(FUN(X-X3*H)+FUN(X+X3*H)))                               00001550
C      3 +W4*(FUN(X-X4*H)+FUN(X+X4*H))) )                             00001560
C      00001570
C      INITIALIZE                                                       00001580
C      00001590
C      IF(ICALL.NE.0)CALL ERRCHK(-71,71H*****GAUS8 CALLED RECURSIVELY. R00001600
C      1ECURSIVE CALLS ARE ILLEGAL IN FORTRAN. )                    00001610
C      ICALL = 1                                                       00001620
C      ANS = 0.0                                                        00001630
C      IERR = 1                                                         00001640
C      CE = 0.0                                                         00001650
C      IF (A.EQ.B) GO TO 35                                             00001660
C      LMX = NLMX                                                       00001670
C      LMN = NLMN                                                       00001680
C      IF (B.EQ.0.0) GO TO 4                                           00001690
C      IF (SIGN(1.0,B)*A.LE.0.0) GO TO 4                               00001700
C      C = ABS(1.0-A/B)                                                00001710
C      IF (C.GT.0.1) GO TO 4                                           00001720

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      IF (C.LE.0.0) GO TO 35
      NIB = 0.5-ALOG(C)/ALOG(2.0)
      LMX = MIN0(NLMX , NBITS-NIB-7)
      IF (LMX.LT.1) GO TO 32
      LMN = MIN0(LMN,LMX)
4    TOL = AMAX1(ABS(ERR) , 2.0**(5-NBITS))/2.0
      IF (ERR.EQ.0.0) TOL = 0.5E-6
      EPS = TOL
      HH(1) = (B-A)/4.0
      AA(1) = A
      LR(1) = 1
      L = 1
      EST = G8(AA(L)+2.0*HH(L) , 2.0*HH(L))
      K = 8
      AREA = ABS(EST)
      EF = 0.5
      MXL = 0
C
C    COMPUTE REFINED ESTIMATES, ESTIMATE THE ERROR, ETC.
C
5    GL = G8(AA(L)+HH(L) , HH(L))
      GR(L) = G8(AA(L)+3.0*HH(L) , HH(L))
      K = K+16
      AREA = AREA+(ABS(GL)+ABS(GR(L))-ABS(EST))
C    IF (L.LT.LMN) GO TO 11
      GLR = GL+GR(L)
      EF = ABS(EST-GLR)*EF
      AE = AMAX1(EPS*AREA , TOL*ABS(GLR))
      IF (EE-AE) 8 , 8 , 10
7    MXL = 1
8    CE = CE + (EST-GLR)
      IF (LR(L)) 15 , 15 , 20
C
C    CONSIDER THE LEFT HALF OF THIS LEVEL
C
10   IF (K.GT.KMX) LMX = KML
      IF (L.GE.LMX) GO TO 7
11   L = L+1
      EPS = EPS*0.5
      EF = EF/SQ2
      HH(L) = HH(L-1)*0.5
      LR(L) = -1
      AA(L) = AA(L-1)
      EST = GL
      GO TO 5
C
C    PROCEED TO RIGHT HALF AT THIS LEVEL
C
15   VL(L) = GLR
16   EST = GR(L-1)
      LR(L) = 1
      AA(L) = AA(L)+4.0*HH(L)
      GO TO 5
C
C    RETURN ONE LEVEL
C
20   VR = GLR
22   IF (L.LE.1) GO TO 30
      L = L-1
      EPS = EPS*2.0
      EF = EF*SQ2
      IF (LR(L)) 24 , 24 , 26
24   VL(L) = VL(L+1)+VR
      GO TO 16
26   VR = VL(L+1)+VR
      GO TO 22
C
C    EXIT
C
30   ANS = VR
      IF ((MXL.EQ.0).OR.(ABS(CE).LE.2.0*TOL*AREA)) GO TO 35
      IERR = 2
      CALL ERRCHK(51,51HIN GAUS8 , ANS IS PROBABLY INSUFFICIENTLY ACCURATE
1TE.)
      GO TO 35
32   IERR = -1

```

```

00001730
00001740
00001750
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00001770
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00001790
00001800
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00002110
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00002190
00002200
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00002230
00002240
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00002390
00002400
00002410
00002420
00002430
00002440
00002450
00002460
00002470
00002480

```

```

CALL ONECHK(-70,70)THE FOLLOWING TEMPORARY INFORMATIVE DIAGNOSTIC 00002490
+ WILL APPEAR ONLY ONCE. ) 00002500
CALLONECHK(-102,102)HIN GAUS8 , A AND B ARE TOO NEARLY EQUAL TO ALL 00002510
LOW NORMAL INTEGRATION. ANS IS SET TO ZERO, AND IERR=-1.) 00002520
35 ICALL = 0 00002530
IF (ERR.LT.0.0) ERR = CE 00002540
RETURN 00002550
END 00002560

SUBROUTINE ERRCHK(NCHARS,NARRAY) 00002570
C 00002580
C SANDIA MATHEMATICAL PROGRAM LIBRARY 00002590
C APPLIED MATHEMATICS DIVISION 2642 00002600
C SANDIA LABORATORIES 00002610
C ALBUQUERQUE, NEW MEXICO 87115 00002620
C 00002630
C SIMPLIFIED VERSION FOR STAND-ALONE USE. APRIL 1977 00002640
C 00002650
C ABSTRACT 00002660
C THE ROUTINES ERRCHK, ERXSET, AND ERRGET TOGETHER PROVIDE 00002670
C A UNIFORM METHOD WITH SEVERAL OPTIONS FOR THE PROCESSING 00002680
C OF DIAGNOSTICS AND WARNING MESSAGES WHICH ORIGINATE 00002690
C IN THE MATHEMATICAL PROGRAM LIBRARY ROUTINES. 00002700
C ERRCHK IS THE CENTRAL ROUTINE, WHICH ACTUALLY PROCESSES 00002710
C MESSAGES. 00002720
C 00002730
C DESCRIPTION OF ARGUMENTS 00002740
C NCHARS - NUMBER OF CHARACTERS IN HOLLERITH MESSAGE. 00002750
C IF NCHARS IS NEGATED, ERRCHK WILL UNCONDITIONALLY 00002760
C PRINT THE MESSAGE AND STOP EXECUTION. OTHERWISE, 00002770
C THE BEHAVIOR OF ERRCHK MAY BE CONTROLLED BY 00002780
C AN APPROPRIATE CALL TO ERXSET. 00002790
C NARRAY - NAME OF ARRAY OR VARIABLE CONTAINING THE MESSAGE, 00002800
C OR ELSE A LITERAL HOLLERITH CONSTANT CONTAINING 00002810
C THE MESSAGE. BY CONVENTION, ALL MESSAGES SHOULD 00002820
C BEGIN WITH *IN SUBNAM, ...*, WHERE SUBNAM IS THE 00002830
C NAME OF THE ROUTINE CALLING ERRCHK. 00002840
C 00002850
C EXAMPLES 00002860
C 1. TO ALLOW CONTROL BY CALLING ERXSET, USE 00002870
C CALL ERRCHK(30,30)HIN QUAD, INVALID VALUE OF ERR.) 00002880
C 2. TO UNCONDITIONALLY PRINT A MESSAGE AND STOP EXECUTION, USE 00002890
C CALL ERRCHK(-30,30)HIN QUAD, INVALID VALUE OF ERR.) 00002900
C 00002910
C 00002920
C 00002930
C ERRCHK USES SUBROUTINES ERRGET, ERRPRT, ERXSET, ERSTGT 00002940
C COMPILE DECKS ERRCHK 00002950
C 00002960
C DIMENSION NARRAY(14) 00002970
C 00002980
C IOUT=6 00002990
C CALL ERRGET(NF,NT) 00003000
C IF ERRCHK WAS CALLED WITH NEGATIVE CHARACTER COUNT, SET FATAL FLAG 00003010
C IF (NCHARS.LT.0) NF = -1 00003020
C IF MESSAGES ARE TO BE SUPPRESSED, RETURN 00003030
C IF (NF.EQ.0) RETURN 00003040
C IF CHARACTER COUNT IS INVALID, STOP 00003050
C IF (NCHARS.EQ.0) PRINT 5 00003060
C IF (NCHARS .EQ. 0) WRITE (IOUT,5) 00003070
5 FORMAT(/31H ERRCHK WAS CALLED INCORRECTLY.) 00003080
C IF (NCHARS.EQ.0) STOP 00003090
C PRINT MESSAGE 00003100
C CALL ERRPRT(IABS(NCHARS),NARRAY) 00003110
C IF LAST MESSAGE, SAY SO 00003120
C IF (NF.EQ.1) PRINT 10 00003130
C IF (NF .EQ. 1) WRITE (IOUT,10) 00003140
10 FORMAT (30H ERRCHK MESSAGE LIMIT REACHED.) 00003150
C PRINT TRACE-BACK IF ASKED TO 00003160
C IF ((NT.GT.0).OR.(NF.LT.0)) CALL SYSTEM ROUTINE FOR TRACEBACK 00003170
C DECREMENT MESSAGE COUNT 00003180
C IF (NF.GT.0) NF = NF-1 00003190
C CALL ERXSET(NF,NT) 00003200
C IF ALL IS WELL, RETURN 00003210
C IF (NF.GE.0) RETURN 00003220

```

```

C   IF THIS MESSAGE IS SUPPRESSABLE BY AN ERXSET CALL,          00003230
C   THEN EXPLAIN ERXSET USAGE.                                  00003240
C   IF (NCHARS.GT.0) PRINT 15                                   00003250
C   IF (NCHARS .GT. 0) WRITE (IOUT,15)                         00003260
15  FORMAT (/13H *** NOTE ***                                   00003270
      1/53H TO MAKE THE ERROR MESSAGE PRINTED ABOVE BE NONFATAL, 00003280
      2/39H OR TO SUPPRESS THE MESSAGE COMPLETELY,              00003290
      3/37H INSERT AN APPROPRIATE CALL TO ERXSET                00003300
      4,30H AT THE START OF YOUR PROGRAM.                       00003310
      5/62H FOR EXAMPLE, TO PRINT UP TO 10 NONFATAL WARNING MESSAGES, USE 00003320
      6/27H          CALL ERXSET(10,0) )                         00003330
C   PRINT 20                                                    00003340
C   WRITE (IOUT,20)                                             00003350
20  FORMAT (/28H PROGRAM ABORT DUE TO ERROR.)                   00003360
C   STOP                                                         00003370
C   END                                                           00003380

SUBROUTINE ONECHK(NCHARS,NARRAY)                                00003390
C   ABSTRACT                                                    00003400
C   ONECHK IS A COMPANION ROUTINE OF ERRCHK. IT IS CALLED      00003410
C   JUST LIKE ERRCHK, AND MESSAGES FROM IT MAY BE SUPPRESSED 00003420
C   BY AN APPROPRIATE CALL TO ERXSET. IT DIFFERS FROM ERRCHK 00003430
C   IN THAT EACH CALL TO ONECHK WILL PRODUCE NO MORE THAN ONE 00003440
C   PRINTED MESSAGE, REGARDLESS OF HOW MANY TIMES THAT CALL IS 00003450
C   EXECUTED, AND ONECHK NEVER TERMINATES EXECUTION.          00003460
C   ITS PURPOSE IS TO PROVIDE ONE-TIME-ONLY INFORMATIVE        00003470
C   DIAGNOSTICS.                                                00003480
C   DESCRIPTION OF ARGUMENTS                                     00003490
C   NCHARS - NUMBER OF CHARACTERS IN THE MESSAGE.              00003500
C   IF NEGATED, THE MESSAGE WILL BE PRINTED (ONCE) EVEN       00003510
C   IF NFATAL HAS BEEN SET TO 0 (SEE ERXSET).                  00003520
C   NARRAY - SAME AS IN ERRCHK                                  00003530
C   ONECHK USES SUBROUTINES ERRGET, ERRPRT, ERXSET, ERSTGT     00003540
C   COMPILER DECKS ERRCHK                                       00003550
C   DIMENSION NARRAY(14)                                        00003560
C   DATA NFLAG/4H.$,*/                                         00003570
C   IF (NARRAY(1).EQ.NFLAG) RETURN                              00003580
C   CALL ERRGET(NF,NT)                                          00003590
C   IF ((NF.EQ.0).AND.(NCHARS.GT.0)) RETURN                    00003600
C   CALL ERRPRT (59,59H THE FOLLOWING INFORMATIVE DIAGNOSTIC WILL APPEAR 00003610
1R ONLY ONCE.)                                                00003620
C   CALL ERRPRT(LABS(NCHARS),NARRAY)                            00003630
C   IF (NF.GT.0) NF = NF-1                                     00003640
C   CALL ERXSET(NF,NT)                                          00003650
C   NARRAY(1) = NFLAG                                          00003660
C   RETURN                                                       00003670
C   END                                                           00003680

SUBROUTINE ERRPRT(NCHARS,NARRAY)                                00003750
C   UTILITY ROUTINE TO SIMPLY PRINT THE HOLLERITH MESSAGE IN NARRAY, 00003760
C   WHOSE LENGTH IS NCHARS CHARACTERS.                          00003770
C   DIMENSION NARRAY(14)                                        00003780
C   NOTE - NCH MUST BE THE NUMBER OF HOLLERITH CHARACTERS STORED 00003790
C   PER WORD. IF NCH IS CHANGED, FORMAT 1 MUST ALSO BE        00003800
C   CHANGED CORRESPONDINGLY.                                   00003810
C   IOUT=6                                                       00003820
C   NCH = 10                                                     00003830
C   FOR LINE PRINTERS, USE                                      00003840
1  FORMAT (1X,13A10)                                           00003850
C   FOR DATA TERMINALS, USE                                    00003860
1  FORMAT (1X,7A10)                                             00003870
C   NWORDS = (NCHARS+NCH-1)/NCH                                  00003880
C   PRINT 1,(NARRAY(I),I=1,NWORDS)                             00003890

```

```

WRITE (IOUT,1) (NARRAY(I),I=1,NWORDS)      00003940
RETURN                                       00003950
END                                           00003960

SUBROUTINE ERXSET(NFATAL,NTRACE)             00003970
C
C ABSTRACT                                    00003980
C ERXSET IS A COMPANION ROUTINE TO SUBROUTINE ERRCHK. 00004000
C ERXSET ASSIGNS THE VALUES OF NFATAL AND NTRACE RESPECTIVELY 00004010
C TO NF AND NT IN COMMON BLOCK MLBLK0 THEREBY SPECIFYING THE 00004020
C STATE OF THE OPTIONS WHICH CONTROL THE EXECUTION OF ERRCHK. 00004030
C DESCRIPTION OF ARGUMENTS                   00004040
C BOTH ARGUMENTS ARE INPUT ARGUMENTS OF DATA TYPE INTEGER. 00004050
C NFATAL - IS A FATAL-ERROR / MESSAGE-LIMIT FLAG. A NEGATIVE 00004070
C VALUE DENOTES THAT DETECTED DIFFICULTIES ARE TO BE 00004080
C TREATED AS FATAL ERRORS. NONNEGATIVE MEANS NONFATAL. 00004090
C A NONNEGATIVE VALUE IS THE MAXIMUM NUMBER OF NONFATAL 00004100
C WARNING MESSAGES WHICH WILL BE PRINTED BY ERRCHK, 00004110
C AFTER WHICH NONFATAL MESSAGES WILL NOT BE PRINTED. 00004120
C (DEFAULT VALUE IS -1.) 00004130
C NTRACE - .GE.1 WILL CAUSE A TRACE-BACK TO BE GIVEN, 00004140
C IF THIS FEATURE IS IMPLEMENTED ON THIS SYSTEM. 00004150
C .LE.0 WILL SUPPRESS ANY TRACE-BACK, EXCEPT FOR 00004160
C CASES WHEN EXECUTION IS TERMINATED. 00004170
C (DEFAULT VALUE IS 0.) 00004180
C
C *NOTE* -- SOME CALLS TO ERRCHK WILL CAUSE UNCONDITIONAL 00004190
C TERMINATION OF EXECUTION. ERXSET HAS NO EFFECT ON SUCH CALLS. 00004210
C
C EXAMPLES                                    00004220
C 1. TO PRINT UP TO 100 MESSAGES AS NONFATAL WARNINGS USE 00004230
C CALL ERXSET(100,0) 00004240
C 2. TO SUPPRESS ALL MATHLIB WARNING MESSAGES USE 00004250
C CALL ERXSET(0,0) 00004260
C
C ERXSET USES SUBROUTINES ERSTGT 00004270
C COMPILER DECKS ERRCHK 00004280
C
C CALL ERSTGT(0,NFATAL,NTRACE) 00004290
C RETURN 00004300
C END 00004310

SUBROUTINE ERRGET(NFATAL,NTRACE)             00004320
C
C ABSTRACT                                    00004330
C ERRGET IS A COMPANION ROUTINE TO SUBROUTINE ERRCHK. 00004340
C ERRGET ASSIGNS TO NFATAL AND NTRACE RESPECTIVELY THE VALUES 00004350
C OF NF AND NT IN COMMON BLOCK MLBLK0 THEREBY ASCERTAINING THE 00004360
C STATE OF THE OPTIONS WHICH CONTROL THE EXECUTION OF ERRCHK. 00004370
C DESCRIPTION OF ARGUMENTS                   00004380
C DESCRIPTION OF ARGUMENTS                   00004390
C BOTH ARGUMENTS ARE OUTPUT ARGUMENTS OF DATA TYPE INTEGER. 00004400
C NFATAL - CURRENT VALUE OF NF (SEE DESCRIPTION OF ERXSET.) 00004410
C NTRACE - CURRENT VALUE OF NT (SEE DESCRIPTION OF ERXSET.) 00004420
C
C CALL ERSTGT(1,NFATAL,NTRACE) 00004430
C RETURN 00004440
C END 00004450

SUBROUTINE ERSTGT(K,NFATAL,NTRACE)           00004460
C
C THIS ROUTINE IS A SLAVE TO ERRGET AND ERRSET WHICH KEEPS 00004470
C THE FLAGS AS LOCAL VARIABLES. 00004480
C
C *** IF LOCAL VARIABLES ARE NOT NORMALLY RETAINED BETWEEN 00004490
C CALLS ON THIS SYSTEM, THE VARIABLES LNF AND LNT CAN BE 00004500
C PLACED IN A COMMON BLOCK AND PRESET TO THE FOLLOWING 00004510

```

COLLECTED ALGORITHMS (cont.)**556-P14- 0**

```
C      VALUES IN THE MAIN PROGRAM.  
C  
      DATA LNF/-1/,LNT/0/  
      IF (K.LE.0) LNF = NFATAL  
      IF (K.LE.0) LNT = NTRACE  
      IF (K.GT.0) NFATAL = LNF  
      IF (K.GT.0) NTRACE = LNT  
      RETURN  
      END
```

```
00004610  
00004620  
00004630  
00004640  
00004650  
00004660  
00004670  
00004680  
00004690
```


ALGORITHM 557

PAGP, A Partitioning Algorithm for (Linear) Goal Programming Problems [H]

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Key Words and Phrases: goal program, multiple objective optimization, constraint partitioning, simplex method
CR Categories: 5.41
Language: Fortran

DESCRIPTION

This algorithm is a Fortran implementation of the procedures developed in [1].

REFERENCE

1. ARTHUR, J.L., AND RAVINDRAN, A. PAGP, A partitioning algorithm for (linear) goal programming problems. *ACM Trans. Math. Softw.* 6, 3 (Sept. 1980), 378-386.

ALGORITHM

```

C      PROGRAM MAIN (INPUT,OUTPUT,TAPE5=INPUT,TAPE6=OUTPUT)          A  10
C      *****
C      ****
C      ****
C      **** PAGP (THE PARTITIONING ALGORITHM FOR GOAL PROGRAMMING) IS
C      **** DESIGNED TO SOLVE THE LINEAR GOAL PROGRAMMING PROBLEM.  THE
C      **** ALGORITHM PARTITIONS THE GOAL CONSTRAINTS OF THE PROBLEM
C      **** ACCORDING TO THE HIGHEST PRIORITY ASSIGNED TO EITHER DEVIATIONAL
C      **** VARIABLE (D- OR D+) IN EACH GOAL.  THE ALGORITHM IS THUS ABLE TO
C      **** SOLVE A SEQUENCE OF SMALLER PROBLEMS IN ORDER TO FIND A SOLUTION
C      **** TO THE ORIGINAL PROBLEM.
C      ****
C      ****
C      *****
C      ****
C      ****
C      **** MUCH OF THE NOTATION AND STRUCTURE OF THE PAGP CODE IS TAKEN
C      **** (WITH THE AUTHOR'S PERMISSION) FROM THE LINEAR GOAL PROGRAMMING

```

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COMMON /CHNG/ NCON(60,10),NTOF(10)	A	50
INTEGER ALTST	A	60
C	A	80
C **** READ IN PROBLEM DATA	A	90
C ****		
C **** NPRIT=THE TOTAL NUMBER OF PRIORITIES		
C ****		
C **** NVAR=THE TOTAL NUMBER OF DECISION VARIABLES(INCLUDING SLACK AND		
C **** SURPLUS VARIABLES BUT EXCLUDING ARTIFICIAL VARIABLES FOR THE		
C **** REAL CONSTRAINTS)		
C ****		
C **** NRCON=THE NUMBER OF REAL CONSTRAINTS		
C	A	100
READ (5,120) NPRIT,NVAR,NRCON	A	110
READ (5,121) (NC(NP),NP=1,NPRIT)	A	120
DO 101 NP=1,NPRIT	A	130
IF (NC(NP).EQ.0) GO TO 101	A	140
NCTMP=NC(NP)		
READ (5,122) (NCON(N,NP),N=1,NCTMP)		
101 CONTINUE	A	160
READ (5,121) (NTOF(NP),NP=1,NPRIT)	A	170
C	A	180
C **** INITIALIZE SUBPROBLEM DIMENSIONS AND COLUMN INDICATORS.	A	190
C ****		
C **** NCOLI=THE NUMBER OF COLUMNS IN THE CURRENT WORKING TABLEAU		
C ****		
C **** NROWI=THE NUMBER OF ROWS IN THE CURRENT WORKING TABLEAU		
C ****		
C **** NPRIC=THE PRIORITY CURRENTLY BEING OPTIMIZED		
C ****		
C **** ZERO THE TE, TL, TT, AND TI ARRAYS.	A	200
C	A	210
NCOLI=0	A	220
NROWI=0	A	230
NPRIC=0	A	240
DO 104 NCR=1,125	A	250
IND(NCR)=1	A	260
DO 102 NR=1,60	A	270
102 TE(NR,NCR)=0.	A	280
DO 103 NP=1,10	A	290
TI(NP,NCR)=0.	A	300
103 TT(NP,NCR)=0.	A	310
104 CONTINUE	A	320
DO 105 NR=1,60	A	330
DO 105 NP=1,10	A	340
105 TL(NR,NP)=0.	A	350
C	A	360
C **** CHECK FOR REAL CONSTRAINTS.	A	370
C	A	380
IF (NRCON.EQ.0) GO TO 106	A	390
CALL PHSE1	A	400
IF (NDVR.LE.0) GO TO 116	A	410
IF (W.GT.0.) GO TO 117	A	420
C	A	430
C **** THE PARTITIONING ALGORITHM BEGINS.	A	440
C	A	450
106 NPRIC=NPRIC+1	A	460
IF (NPRIC.EQ.1.AND.NRCON.EQ.0) GO TO 107	A	470
GO TO 108	A	480
107 CALL READ1	A	490
GO TO 109	A	500
108 CALL READ2	A	510
109 CALL CINDX	A	520
CALL TEST (NEVC,NDVR)	A	530
C	A	540
C **** IF NEVC IS LESS THAN ZERO, THE SUBPROBLEM IS OPTIMIZED.	A	550
C	A	560
IF (NEVC.LE.0) GO TO 110	A	570
C	A	580
C **** IF NDVR IS LESS THAN ZERO, NO MINIMUM POSITIVE RATIO WAS FOUND.	A	590
C	A	600
IF (NDVR.LE.0) GO TO 116	A	610
CALL PERM (NEVC,NDVR)	A	620
GO TO 109	A	630
C	A	640
C **** IF THERE ARE NO MORE PRIORITIES, TOTAL PROBLEM IS OPTIMIZED.	A	650

```

C **** PRINT THE OPTIMAL SOLUTION. A 660
C A 670
C 110 IF (NPRIC.EQ.NPRIT) GO TO 115 A 680
C A 690
C **** SINCE THERE ARE MORE PRIORITIES, MOVE ON TO THE NEXT SUBPROBLEM A 700
C **** IF THERE ARE ALTERNATE SOLUTIONS. FIRST, ELIMINATE THOSE A 710
C **** COLUMNS WHICH CAN NOT ENTER THE BASIS. IF THERE ARE NO A 720
C **** ALTERNATE SOLUTIONS, PRINT THE UNIQUE OPTIMAL SOLUTION. A 730
C A 740
C ALTST=0 A 750
C DO 112 NCR=1,NCOLI A 760
C IF (IND(NCR).EQ.0) GO TO 112 A 770
C IF (TI(NPRIC,NCR).GT.0.) GO TO 112 A 780
C DO 111 NR=1,NROWI A 790
C IF (JROW(NR,1).EQ.JCOL(NCR,1).AND.JROW(NR,2).EQ.JCOL(NCR,2)) A 800
C 1 GO TO 112 A 810
C 111 CONTINUE A 820
C ALTST=1 A 830
C 112 CONTINUE A 840
C A 850
C **** IF ALTST=1, THERE ARE ALTERNATE SOLUTIONS. A 860
C A 870
C IF (ALTST.EQ.1) GO TO 113 A 880
C GO TO 115 A 890
C A 900
C **** ELIMINATE THOSE COLUMNS WITH A POSITIVE RELATIVE COST AT A 910
C **** PRIORITY NPRIC. A 920
C A 930
C 113 DO 114 NCR=1,NCOLI A 940
C 114 IF (TI(NPRIC,NCR).GT.0.) IND(NCR)=0 A 950
C GO TO 106 A 960
C
C **** THE OPTIMIZATION IS OVER. PRINT OUT THE FINAL SOLUTION.
C
C 115 CALL POUT A 970
C GO TO 119 A 980
C 116 WRITE (6,123) NPRIC A 990
C GO TO 119 A 1000
C 117 WRITE (6,124) W A 1010
C WRITE (6,125) A 1020
C DO 118 NR=1,NROWI A 1030
C WRITE (6,126) JROW(NR,1),JROW(NR,2),TB(NR) A 1040
C 118 CONTINUE A 1050
C 119 STOP
C
C 120 FORMAT (3I5) A 1120
C 121 FORMAT (10I5) A 1130
C 122 FORMAT (16I5) A 1150
C 123 FORMAT (/ 40H THE PROGRAM TERMINATED ON SUBPROBLEM ,I4, 42H NO A 1170
C 1 MINIMUM POSITIVE RATIO COULD BE FOUND)
C 124 FORMAT (/ 65H THE PROGRAM TERMINATED IN PHASE 1 WITH OBJECTIVE F A 1190
C LUNCTION VALUE,F15.4)
C 125 FORMAT (/ 55H THE OPTIMAL SOLUTION TO THE PHASE 1 PROBLEM IS
C 1 // 6H TYPE,2X, 3H SUB,8X, 5H VALUE)
C 126 FORMAT (2I5,F15.4) A 1220
C A 1260
C END A 1270

SUBROUTINE PHSE1 B 10

C
C **** SUBROUTINE PHSE1 READS IN ANY REAL CONSTRAINTS AND PERFORMS A
C **** PHASE 1 SIMPLEX PROCEDURE IN ORDER TO FIND AN INITIAL BASIC
C **** FEASIBLE SOLUTION.
C
C COMMON TT(10,125),TB(60),TE(60,125),TL(60,10),TA(10),TI(10,125),JC B 20
C IOL(125,2),NCOLI,NROWI,NPRIC,NC(10),JROW(60,2),NVAR,NPRIT,IND(125) B 30
C COMMON /PHASE1/ W,NRCON,NDVR B 40
C DIMENSION C(125), CR(125), CB(60) B 50
C B 60
C **** SET COLUMN AND ROW HEADINGS B 70
C B 80
C DO 101 NV=1,NVAR B 90
C JCOL(NV,1)=2 B 100
C 101 JCOL(NV,2)=NV B 110
C DO 102 NR=1,NRCON B 120

```

	JROW(NR,1)=1	B 130
	JROW(NR,2)=NR	B 140
	NAR=NVAR+NR	B 150
	JCOL(NAR,1)=1	B 160
	102 JCOL(NAR,2)=NR	B 170
C		B 180
C	**** READ IN COEFFICIENTS AND RHS OF REAL CONSTRAINTS	B 190
C		B 200
	DO 103 NR=1,NRCON	B 210
	READ (5,118) TB(NR),(TE(NR,NV),NV=1,NVAR)	B 220
	103 CONTINUE	B 230
C		B 240
C	**** PUT IDENTITY MATRIX IN FOR ARTIFICIAL VARIABLES	B 250
C		B 260
	DO 104 NR=1,NRCON	B 270
	NAR=NVAR+NR	B 280
	104 TE(NR,NAR)=1.	B 290
C		B 300
C	**** SET C(J)=0 FOR ALL DECISION VARIABLES AND C(J)=1 FOR ALL	B 310
C	**** ARTIFICIAL VARIABLES	B 320
C		B 330
	DO 105 NV=1,NVAR	B 340
	105 C(NV)=0.	B 350
	DO 106 NR=1,NRCON	B 360
	CB(NR)=1.	B 370
	NAR=NVAR+NR	B 380
	106 C(NAR)=1.	B 390
C		B 400
C	**** CALCULATE RELATIVE COST COEFFICIENTS CR(.)	B 410
C		B 420
	NCOL=NVAR+NRCON	B 430
	107 DO 108 NV=1,NCOL	B 440
	CR(NV)=C(NV)	B 450
	DO 108 NR=1,NRCON	B 460
	108 CR(NV)=CR(NV)-CB(NR)*TE(NR,NV)	B 470
C		B 480
C	**** CHECK FOR OPTIMALITY	B 490
C		B 500
	VEVC=0.	B 510
	NEVC=0	B 520
	DO 109 NCO=1,NCOL	B 530
	NV=NCO	B 540
	IF (CR(NV).GE.0.) GO TO 109	B 550
	IF (CR(NV).GE.VEVC) GO TO 109	B 560
	VEVC=CR(NV)	B 570
	NEVC=NV	B 580
	109 CONTINUE	B 590
C		B 600
C	**** IF NEVC=0, PHASE 1 IS OPTIMIZED. CALCULATE OBJECTIVE FUNCTION.	B 610
C		B 620
	IF (NEVC.EQ.0) GO TO 115	B 630
C		B 640
C	**** DETERMINE DEPARTING VARIABLES ROW.	B 650
C		B 660
	NDVR=0	B 670
	VDVR=10.0E+20	B 680
	DO 111 NRI=1,NRCON	B 690
	NR=NRI	B 700
	IF (TE(NR,NEVC).LE.0.) GO TO 111	B 710
	V=TB(NR)/TE(NR,NEVC)	B 720
	IF (NDVR.EQ.0) GO TO 110	B 730
	IF (V-VDVR) 110,110,111	B 740
	110 VDVR=V	B 750
	NDVR=NR	B 760
	111 CONTINUE	B 770
C		B 780
C	**** IF NDVR=0, MINIMUM RATIO RULE FAILED. RETURN.	B 790
C		B 800
	IF (NDVR.EQ.0) RETURN	B 810
C		B 820
C	**** PERFORM THE PIVOT. REPLACE HEADINGS AND COST COEFFICIENT.	B 830
C		B 840
	JROW(NDVR,1)=JCOL(NEVC,1)	B 850
	JROW(NDVR,2)=JCOL(NEVC,2)	B 860
	CB(NDVR)=C(NEVC)	B 870
C		B 880

```

C **** COMPUTE NEW TE ARRAY
C
  PIV=TE(NDVR,NEVC)
  PIB=TB(NDVR)
  DO 113 NR=1,NRCON
    IF (NR.EQ.NDVR) GO TO 113
    IF (ABS(TE(NR,NEVC)).LE.0.0005) GO TO 113
    PIX=TE(NR,NEVC)/PIV
    TB(NR)=FIX(TB(NR)-PIX*PIB)
    DO 112 NV=1,NCOL
      112 TE(NR,NV)=FIX(TE(NR,NV)-TE(NDVR,NV)*PIX)
  113 CONTINUE
    TB(NDVR)=FIX(PIB/PIV)
    DO 114 NV=1,NCOL
      114 TE(NDVR,NV)=FIX(TE(NDVR,NV)/PIV)
C
C **** END OF PIVOT OPERATIONS. PROCEED TO NEXT ITERATION.
C
  GO TO 107
C
C **** CALCULATE W, THE PHASE-1 OBJECTIVE FUNCTION.
C
  115 W=0.
  DO 116 NR=1,NRCON
    116 W=W+TB(NR)*CB(NR)
C
C **** INITIALIZE THOSE PORTIONS OF THE TABLEAU ASSIGNED TO THE
C **** ARTIFICIAL VARIABLES.
C
  DO 117 NR=1,NRCON
    DO 117 NV=1,NCOL
      IF (NV.LE.NVAR) GO TO 117
      TE(NR,NV)=0.
  117 CONTINUE
C
C **** UPDATE NCOLI AND NROWI PARAMETERS.
C
  NROWI=NRCON
  NCOLI=NVAR
  RETURN
C
  118 FORMAT (8F10.0)
C
  END

```

B 890
B 900
B 910
B 920
B 930
B 940
B 950
B 960
B 970
B 980
B 990
B 1000
B 1010
B 1020
B 1030
B 1040
B 1050
B 1060
B 1070
B 1080
B 1100
B 1110
B 1120
B 1130
B 1140
B 1150
B 1160
B 1170
B 1180
B 1190
B 1200
B 1210
B 1220
B 1230
B 1240
B 1250
B 1260
B 1270
B 128
B 1290
B 1300
B 1310
B 1320

C 10

```

SUBROUTINE READ1
C
C **** SUBROUTINE READ1 READS IN THE GOAL CONSTRAINTS AND OBJECTIVE
C **** FUNCTION TERMS ASSIGNED TO PRIORITY ONE.
C **** SUBROUTINE READ1 IS NOT USED IF REAL CONSTRAINTS ARE PRESENT.
C
  COMMON TT(10,125),TB(60),TE(60,125),TL(60,10),TA(10),TI(10,125),JC
  IOL(125,2),NCOLI,NROWI,NPRIC,NC(10),JROW(60,2),NVAR,NPRIT,IND(125)
  COMMON /CHNG/ NCON(60,10),NTOF(10)
C
C **** SET COLUMN AND ROW HEADINGS.
C
  DO 101 NV=1,NVAR
    JCOL(NV,1)=2
  101 JCOL(NV,2)=NV
    NCL1=NC(1)
    DO 102 NCR=1,NC11
      NC1=NVAR+2*NCR-1
      NC2=NVAR+2*NCR
      JCOL(NC1,1)=4
      JCOL(NC1,2)=NCON(NCR,1)
      JCOL(NC2,1)=3
      JCOL(NC2,2)=NCON(NCR,1)
      JROW(NCR,1)=4
  102 JROW(NCR,2)=NCON(NCR,1)
C
C **** READ IN THE GOAL CONSTRAINTS ASSIGNED TO PRIORITY 1.
C
  NC1=NC(1)
  DO 103 NCR=1,NC1

```

C 20
C 30
C 40
C 50
C 60
C 70
C 80
C 90
C 100
C 120
C 130
C 140
C 150
C 160
C 170
C 180
C 190
C 200
C 210
C 220

```

      NV1=NVAR+2*NCR-1          C 240
      NV2=NVAR+2*NCR          C 250
      READ (5,105) TB(NCR),(TE(NCR,NV),NV=1,NVAR) C 260
C                                     C 270
C **** PUT +1 IN FOR D- AND -1 IN FOR D+. C 280
C                                     C 290
      TE(NCR,NV1)=1.          C 300
      TE(NCR,NV2)=-1.        C 310
103 CONTINUE                  C 320
      NCOLI=NV2                C 330
      NROWI=NC(1)              C 340
C                                     C 350
C **** READ IN THE OBJECTIVE FUNCTION TERMS FOR PRIORITY 1. C 360
C                                     C 370
      NT1=NTOF(1)
      DO 104 NT=1,NT1
        READ (5,106) ISUB,IYPE,WGHT C 390
        CALL PLACE (ISUB,IYPE,WGHT) C 400
104 CONTINUE                  C 410
      RETURN                   C 420
C                                     C 430
105 FORMAT (8F10.0)          C 440
106 FORMAT (2I5,F10.0)      C 450
C                                     C 460
      END                       C 470

      SUBROUTINE READ2          D 10
C
C **** SUBROUTINE READ2 READS IN THE GOAL CONSTRAINTS AND OBJECTIVE
C **** FUNCTION TERMS ASSIGNED TO PRIORITY NPRIC.
C **** SUBROUTINE READ2 IS ALSO USED TO READ IN THE FIRST PRIORITY GOAL
C **** CONSTRAINTS AND OBJECTIVE FUNCTION TERMS IF REAL CONSTRAINTS ARE
C **** PRESENT.
C
      COMMON TT(10,125),TB(60),TE(60,125),TL(60,10),TA(10),TI(10,125),JC D 20
10L(125,2),NCOLI,NROWI,NPRIC,NC(10),JROW(60,2),NVAR,NPRIT,IND(125) D 30
      COMMON /CHNG/ NCON(60,10),NTOF(10) D 40
      IF (NC(NPRIC).EQ.0) GO TO 107 D 50
C                                     D 60
C **** READ IN THE COEFFICIENTS OF THE X'S. D 70
C                                     D 80
      NCTMP=NC(NPRIC)
      DO 106 NRI=1,NCTMP
        NR=NRI+NROWI          D 100
        NC1=NCOLI+2*NRI-1     D 110
        NC2=NCOLI+2*NRI     D 120
        JCOL(NC1,1)=4         D 130
        JCOL(NC1,2)=NCON(NRI,NPRIC) D 140
        JCOL(NC2,1)=3         D 150
        JCOL(NC2,2)=NCON(NRI,NPRIC) D 160
        READ (5,109) TB(NR),(TE(NR,NV),NV=1,NVAR) D 170
        TE(NR,NC1)=1.         D 180
        TE(NR,NC2)=-1.       D 190
C                                     D 200
C **** PERFORM THE ROW REDUCTION. D 210
C                                     D 220
      DO 102 NRC=1,NROWI      D 230
        IF (JROW(NRC,1).NE.2) GO TO 102 D 240
        J=JROW(NRC,2)         D 250
        TB(NR)=TB(NR)-TE(NR,J)*TB(NRC) D 260
        DO 101 NCR=1,NC2     D 270
          IF (NCR.EQ.J) GO TO 101 D 280
          TE(NR,NCR)=TE(NR,NCR)-TE(NR,J)*TE(NRC,NCR) D 290
101 CONTINUE                  D 300
          TE(NR,J)=0.         D 310
102 CONTINUE                  D 320
C                                     D 330
C **** DETERMINE THE DEVIATIONAL VARIABLE TO ENTER THE BASIS. D 340
C                                     D 350
      IF (TB(NR)) 103,105,105 D 360
C                                     D 370
C **** SINCE TB IS LESS THAN ZERO, MULTIPLY THE ROW BY -1 AND ENTER D+ D 380
C **** IN THE BASIS. D 390
C                                     D 400
103 DO 104 NCR=1,NC2         D 410

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104   TE(NR,NCR)=-TE(NR,NCR)           D 420
      TB(NR)=-TB(NR)                   D 430
      JROW(NR,1)=3                       D 440
      JROW(NR,2)=NCON(NRI,NPRIC)        D 450
      GO TO 106                           D 460
C                                           D 470
C **** SINCE TB IS GREATER THAN OR EQUAL TO ZERO ENTER D- IN THE BASIS. D 480
C                                           D 490
105   JROW(NR,1)=4                       D 500
      JROW(NR,2)=NCON(NRI,NPRIC)        D 510
106   CONTINUE                           D 520
C                                           D 530
C **** INCREASE THE PARAMETERS NCOLI AND NROWI. D 540
C                                           D 550
      NCOLI=NC2                           D 560
      NROWI=NR                             D 570
C                                           D 580
C **** READ IN THE OBJECTIVE FUNCTION TERMS FOR PRIORITY NPRIC. D 590
C                                           D 600
107   NTTMP=NTOF(NPRIC)
      DO 108 NT=1,NTTMP
          READ (5,110) ISUB,ITYPE,WGHT D 620
          CALL PLACE (ISUB,ITYPE,WGHT) D 630
108   CONTINUE                           D 640
      RETURN                               D 650
C                                           D 660
109   FORMAT (8F10.0)                     D 670
110   FORMAT (2I5,F10.0)                   D 680
C                                           D 690
      END                                   D 700

      SUBROUTINE PLACE (ISUB,ITYPE,WGHT) E 10
C
C **** SUBROUTINE PLACE PUTS THE OBJECTIVE FUNCTION WEIGHTS FOR THE
C **** DEVIATION VARIABLES AT THE CURRENT PRIORITY LEVEL (NPRIC) IN THE
C **** CORRECT POSITIONS IN THE AUGMENTED TABLEAU.
C ****
C **** ISUB=THE SUBSCRIPT OF THE DEVIATIONAL VARIABLE
C ****
C **** ITYPE=3, IF POSITIVE DEVIATIONAL VARIABLE (D+)
C ****         4, IF NEGATIVE DEVIATIONAL VARIABLE (D-)
C ****
C **** WGHT=THE CARDINAL WEIGHT OF THIS DEVIATIONAL VARIABLE AT THE
C **** CURRENT PRIORITY LEVEL
C
      COMMON TT(10,125),TB(60),TE(60,125),TL(60,10),TA(10),TI(10,125),JC E 20
      IOL(125,2),NCOLI,NROWI,NPRIC,NC(10),JROW(60,2),NVAR,NPRIT,IND(125) E 30
      COMMON /CHNG/ NCON(60,10),NTOF(10) E 40
C                                           E 50
C **** PLACE THE WEIGHT IN THE PROPER COLUMN IN THE TOP STUB. E 60
C                                           E 70
      NCI=NVAR+1                           E 80
      DO 101 NCR=NC1,NCOLI                 E 90
          IF (JCOL(NCR,1).EQ.ITYPE.AND.JCOL(NCR,2).EQ.ISUB) GO TO 102 E 100
101   CONTINUE                             E 110
      102 TT(NPRIC,NCR)=WGHT                E 120
C                                           E 130
C **** PLACE THE WEIGHT IN THE PROPER ROW IN THE LEFT STUB. E 140
C                                           E 150
      DO 103 NR=1,NROWI                    E 160
          IF (JROW(NR,1).EQ.ITYPE.AND.JROW(NR,2).EQ.ISUB) GO TO 104 E 170
103   CONTINUE                             E 180
      GO TO 105                             E 190
104   TL(NR,NPRIC)=WGHT                    E 200
105   CONTINUE                             E 210
      RETURN                               E 220
C                                           E 230
      END                                   E 240

      SUBROUTINE CINDX                       F 10
C
C **** SUBROUTINE CINDX COMPUTES THE RELATIVE COST COEFFICIENTS FOR EACH
C **** VARIABLE IN THE CURRENT TABLEAU(THE TI(. , .) ARRAY) AND THE
C **** OBJECTIVE FUNCTION VALUE(THE TA(.) ARRAY) AT THE CURRENT

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C **** PRIORITY (NPRIC)
C
COMMON TT(10,125),TB(60),TE(60,125),TL(60,10),TA(10),TI(10,125),JC F 20
IOL(125,2),NCOLI,NROWI,NPRIC,NC(10),JROW(60,2),NVAR,NPRIT,IND(125) F 30
C F 40
C **** COMPUTE TA(NPRIC) AND TI(NPRIC,NC) NC=1,...,NCOLI F 50
C F 60
TA(NPRIC)=0. F 70
DO 101 NR=1,NROWI F 80
101 TA(NPRIC)=TA(NPRIC)+TB(NR)*TL(NR,NPRIC) F 90
DO 102 NCR=1,NCOLI F 100
TI(NPRIC,NCR)=TT(NPRIC,NCR) F 110
DO 102 NR=1,NROWI F 120
102 TI(NPRIC,NCR)=TI(NPRIC,NCR)-TE(NR,NCR)*TL(NR,NPRIC) F 130
RETURN F 140
C F 150
END F 160

SUBROUTINE TEST (NEVC,NDVR) G 10
C
C **** SUBROUTINE TEST DETERMINES THE NEXT ENTERING VARIABLE'S COLUMN
C **** (NEVC) AND THE NEXT DEPARTING VARIABLE'S ROW(NDVR). IF NO
C **** FURTHER OPTIMIZATION IS POSSIBLE, THE VALUE NEVC=0 IS RETURNED.
C **** IF NDVR=0 IS RETURNED, NO MINIMUM POSITIVE RATIO COULD BE FOUND
C **** IN THE CURRENT PIVOT OPERATION, I.E., ALL OF THE COEFFICIENTS
C **** TE( . ,NEVC) ARE NONPOSITIVE.
C
COMMON TT(10,125),TB(60),TE(60,125),TL(60,10),TA(10),TI(10,125),JC G 20
IOL(125,2),NCOLI,NROWI,NPRIC,NC(10),JROW(60,2),NVAR,NPRIT,IND(125) G 30
NDVR=0 G 40
NEVC=0 G 50
VEVC=0. G 60
VDVR=10.0E+20 G 70
C G 80
C **** DETERMINE ENTERING VARIABLE'S COLUMN. G 90
C G 100
DO 101 NCR=1,NCOLI G 110
IF (TI(NPRIC,NCR).GE.0.) GO TO 101 G 120
IF (IND(NCR).EQ.0) GO TO 101 G 130
IF (TI(NPRIC,NCR).GE.VEVC) GO TO 101 G 140
NEVC=NCR G 150
VEVC=TI(NPRIC,NCR) G 160
101 CONTINUE G 170
C G 180
C **** IF NEVC=0, SUBPROBLEM NPRIC IS OPTIMIZED. RETURN. G 190
C G 200
IF (NEVC.EQ.0) RETURN G 210
C G 220
C **** DETERMINE DEPARTING VARIABLE'S ROW. G 230
C G 240
DO 105 NR=1,NROWI G 250
IF (TE(NR,NEVC).LE.0.) GO TO 105 G 260
V=TB(NR)/TE(NR,NEVC) G 270
IF (NDVR.EQ.0) GO TO 104 G 280
IF (V-VDVR) 104,102,105 G 290
102 DO 103 NP=1,NPRIC G 300
IF (TL(NR,NP)-TL(NDVR,NP)) 105,103,104 G 310
103 CONTINUE G 320
104 VDVR=V G 330
NDVR=NR G 340
105 CONTINUE G 350
RETURN G 360
C G 370
END G 380

SUBROUTINE PERM (NEVC,NDVR) H 10
C
C **** SUBROUTINE PERM PERFORMS THE PIVOT OPERATION USING THE PIVOT
C **** ELEMENT IN COLUMN NEVC AND ROW NDVR AND COMPUTES THE NEW TABLEAU.
C
COMMON TT(10,125),TB(60),TE(60,125),TL(60,10),TA(10),TI(10,125),JC H 20
IOL(125,2),NCOLI,NROWI,NPRIC,NC(10),JROW(60,2),NVAR,NPRIT,IND(125) H 30
C H 40

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```

C **** REPLACE HEADING FOR ROW NDVR.          H  50
C                                               H  60
      JROW(NDVR,1)=JCOL(NEVC,1)                H  70
      JROW(NDVR,2)=JCOL(NEVC,2)                H  80
C                                               H  90
C **** REPLACE TL VECTOR FOR ROW NDVR          H 100
C                                               H 110
      DO 101 NP=1,NPRIC                         H 120
101  TL(NDVR,NP)=TT(NP,NEVC)                   H 130
C                                               H 140
C **** COMPUTE NEW TE ARRAY.                  H 150
C                                               H 160
      PIV=TE(NDVR,NEVC)                         H 170
      PIB=TB(NDVR)                              H 180
      DO 103 NR=1,NROWI                         H 190
        IF (NR.EQ.NDVR) GO TO 103               H 200
        IF (ABS(TE(NR,NEVC)).LE.0.0005) GO TO 103 H 210
        PIX=TE(NR,NEVC)/PIV                    H 220
        TB(NR)=FIX(TB(NR)-PIX*PIB)              H 230
        DO 102 NCR=1,NCOLI                     H 240
102   TE(NR,NCR)=FIX(TE(NR,NCR)-TE(NDVR,NCR)*PIX) H 250
103  CONTINUE                                  H 260
      TB(NDVR)=FIX(PIB/PIV)                    H 270
      DO 104 NCR=1,NCOLI                       H 280
104  TE(NDVR,NCR)=FIX(TE(NDVR,NCR)/PIV)        H 290
      RETURN                                    H 300
C                                               H 310
      END                                       H 320

      FUNCTION FIX(Z)                            I  10
C
C **** FUNCTION FIX BRINGS FLOATING POINT VALUES THAT ARE WITHIN 1.E-3
C **** OF AN INTEGER TO THAT INTEGER.
C
      FIX=AINT(Z+SIGN(.5,Z))
      IF (ABS(FIX-Z).GT. 1.E-3) FIX=Z
      RETURN                                    I 150
C                                               I 160
      END                                       I 170

      SUBROUTINE POUT                            J  10
C
C **** SUBROUTINE POUT PREPARES AND PRINTS THE SOLUTION INFORMATION.
C
      COMMON TT(10,125),TB(60),TE(60,125),TL(60,10),TA(10),TI(10,125),JC J  20
      LOL(125,2),NCOLI,NROWI,NPRIC,NC(10),JROW(60,2),NVAR,NPRIT,IND(125) J  30
      COMMON /CHNG/ NCON(60,10),NTOF(10)        J  40
      DIMENSION WOUT(125,4), RLHS(60,10), WM(60), WP(60) J  50
      DIMENSION DIFF(60)                       J  60
      WRITE (6,122)                             J  70
      WRITE (6,123) NPRIC,NROWI                 J  80
C                                               J  90
C **** OUTPUT ARRAY IS ZEROED.                J 100
C                                               J 110
      DO 101 I=1,125                             J 120
        DO 101 J=1,4                             J 130
101  WOUT(I,J)=0.                               J 140
C                                               J 150
C **** OUTPUT ARRAY IS FILLED.                J 160
C                                               J 170
      DO 102 NP=1,NPRIC                         J 180
102  WOUT(NP,1)=FIX(TA(NP))                    J 190
      DO 103 NR=1,NROWI                         J 200
        I1=JROW(NR,1)                           J 210
        I2=JROW(NR,2)                           J 220
103  WOUT(I2,I1)=FIX(TB(NR))                   J 230
C                                               J 240
C **** IF ALL PRIORITIES HAVE BEEN INCLUDED, PRINT OPTIMAL SOLUTION. J 250
C **** IF NOT, WE MUST CALCULATE VALUES FOR REMAINING TA'S AND D- AND D+ J 260
C                                               J 270
      IF (NPRIC.GE.NPRIT) GO TO 114            J 280
      NP1=NPRIC+1                               J 290
      DO 113 NP=NP1,NPRIT                       J 300

```

	TA(NP)=0.	J	310
	IF (NC(NP).EQ.0) GO TO 106	J	320
C		J	330
C	**** READ IN THE GOAL CONSTRAINTS ASSIGNED TO PRIORITY NP.	J	340
C		J	350
	NCTMP=NC(NP)		
	DO 105 NCI=1,NCTMP		
	NR=NROWI+NCI	J	370
	READ (5,124) TB(NR),(TE(NR,NV),NV=1,NVAR)	J	380
	RLHS(NCI,NP)=0.	J	390
	DO 104 NV=1,NVAR	J	400
104	RLHS(NCI,NP)=RLHS(NCI,NP)+TE(NR,NV)*WOUT(NV,2)	J	410
	DIFF(NCI)=TB(NR)-RLHS(NCI,NP)	J	420
105	CONTINUE	J	430
C		J	440
C	**** READ THE OBJECTIVE FUNCTION TERMS FOR PRIORITY NP.	J	450
C		J	460
106	NTTMP=NTOF(NP)		
	DO 112 NT=1,NTTMP		
	READ (5,125) ISUB,ITYPE,WGHT	J	480
	IF (NC(NP).EQ.0) GO TO 111	J	490
	NCTMP=NC(NP)		
	DO 110 NCI=1,NCTMP		
	IF (ISUB.NE.NCON(NCI,NP)) GO TO 110	J	510
	IF (DIFF(NCI)) 107,108,109	J	520
107	IF (ITYPE.NE.3) GO TO 110	J	530
	WOUT(ISUB,3)=-DIFF(NCI)	J	540
108	GO TO 110	J	550
109	IF (ITYPE.NE.4) GO TO 110	J	560
	WOUT(ISUB,4)=DIFF(NCI)	J	570
110	CONTINUE	J	580
111	TA(NP)=TA(NP)+WGHT*WOUT(ISUB,ITYPE)	J	590
112	CONTINUE	J	600
	NROWI=NROWI+NC(NP)	J	610
C		J	620
C	**** FILL IN THE OUTPUT VALUE FOR ATTAINMENT OF PRIORITY NP.	J	630
C		J	640
	WOUT(NP,1)=FIX(TA(NP))	J	650
113	CONTINUE	J	660
C		J	670
C	**** PRINT OPTIMAL SOLUTION	J	680
C		J	690
114	WRITE (6,126)	J	700
	WRITE (6,127)	J	710
	DO 115 NV=1,NVAR	J	720
	WRITE (6,128) NV,WOUT(NV,2)	J	730
115	CONTINUE	J	740
	WRITE (6,126)	J	750
	WRITE (6,129)	J	760
	DO 116 NP=1,NPRIT	J	770
	IF (NC(NP).EQ.0) GO TO 116	J	780
	NCTMP=NC(NP)		
	DO 139 NCO=1,NCTMP		
	N=NCON(NCO,NP)	J	800
	WRITE (6,130) NP,N,WOUT(N,3),WOUT(N,4)	J	810
139	CONTINUE		
116	CONTINUE	J	820
	WRITE (6,126)	J	830
	WRITE (6,131)	J	840
	DO 117 NP=1,NPRIT	J	850
	WRITE (6,132) NP,WOUT(NP,1)	J	860
117	CONTINUE	J	870
	WRITE (6,126)	J	880
	WRITE (6,133)	J	890
	WRITE (6,134)	J	900
	I=MAX0(NPRIT,NVAR,NROWI)	J	910
	DO 121 K=1,I	J	920
	IF (K.GT.NPRIT) GO TO 119	J	930
	IF (K.GT.NVAR) GO TO 118	J	940
	WRITE (6,135) K,(WOUT(K,J),J=1,4)	J	950
	GO TO 121	J	960
118	WRITE (6,136) K,WOUT(K,1),(WOUT(K,J),J=3,4)	J	970
	GO TO 121	J	980
119	IF (K.GT.NVAR) GO TO 120	J	990
	WRITE (6,137) K,(WOUT(K,J),J=2,4)	J	1000
	GO TO 121	J	1010

120	WRITE (6,138) K,(WOUT(K,J),J=3,4)	J 1020
121	CONTINUE	J 1030
	WRITE (6,126)	J 1040
	RETURN	J 1050
C		J 1060
122	FORMAT (1H1)	J 1070
123	FORMAT (/ 39H THE OPTIMIZATION ENDED ON SUBPROBLEM ,I5 / 13H T HERE WERE ,I5, 42H CONSTRAINTS IN THE FINAL OPTIMAL TABLEAU.)	J 1090
124	FORMAT (8F10.0)	J 1100
125	FORMAT (2I5,F10.0)	J 1110
126	FORMAT (//120(1H*))	J 1120
127	FORMAT (1H0, 52HTHE OPTIMAL SOLUTION FOR THE DECISION VARIABLES X(1J))	J 1130
128	FORMAT (1H0, 2HX(,I3, 2H)=,F15.4)	J 1150
129	FORMAT (1H0, 25HTHE GOAL ACHIEVEMENTS ARE // 9H PRIORITY,2X, 11H 1GOAL NUMBER,2X, 16HOVER-ACHIEVEMENT,2X, 17HUNDER-ACHIEVEMENT)	J 1170
130	FORMAT (4X,I2,10X,I2,10X,F10.4,10X,F10.4)	J 1180
131	FORMAT (1H0, 29HTHE PRIORITY ACHIEVEMENTS ARE // 9H PRIORITY,8X, 1 11HACHIEVEMENT)	J 1200
132	FORMAT (4X,I2,10X,F10.4)	J 1210
133	FORMAT (1H0, 15H OUTPUT SUMMARY)	J 1220
134	FORMAT (1H0, 9HSUBSCRIPT,11X, 8H A OPT,7X, 8H X OPT,7X, 9H 1 POS DEV,6X, 9H NEG DEV /)	J 1230
135	FORMAT (I8,7X,4F15.4)	J 1250
136	FORMAT (I8,7X,F15.4,15X,2F15.4)	J 1260
137	FORMAT (I8,22X,3F15.4)	J 1270
138	FORMAT (I8,37X,2F15.4)	J 1280
C		J 1290
	END	J 1300

ALGORITHM 558

A Program for the Multifacility Location Problem with Rectilinear Distance by the Minimum-Cut Approach [H]

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University of Ottawa

Key Words and Phrases: multifacility, optimal location, rectilinear distance, minimum cut
CR Categories: 3.57, 5.41
Language: Fortran

DESCRIPTION

This algorithm complements [1], where more details are given. The coding of the algorithm consists of two subroutines, LOCATE and NETFLO. A user's program calls LOCATE, which contains the major coding, and calls NETFLO for solving the minimum-cut network flow subproblems.

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REFERENCE

- CHEUNG, T.Y. Multifacility location problem with rectilinear distance by the minimum cut approach. *ACM Trans. Math. Softw.* 6, 3 (Sept. 1980), 387-390.

ALGORITHM

```

C                                     TP1 0000
C *TEST PROGRAM #1* (DATA INPUT FROM CARDS)           TP1 0001
C THIS IS THE FIRST OF TWO PROGRAMS FOR THE TESTING OF THE   TP1 0002
C SUBROUTINE 'LOCATE'. ALL IT DOES IS TO READ THE DATA NVPT, TP1 0003
C NFPT, VWT AND VFWT FROM CARDS, CALL 'LOCATE', AND PRINT THE TP1 0004
C SOLUTION POST.                                           TP1 0005
C IT FIRST READS AN INTEGER NPRLM FROM A CARD, USING        TP1 0006
C FORMAT(I3), WHERE NPRLM IS THE NUMBER OF PROBLEMS TO BE   TP1 0007
C TESTED IN A SINGLE RUN. THEN, FOR EACH PROBLEM TESTED,   TP1 0008
C A SET OF 1+2*NVPT DATA CARDS WILL BE READ. THE FORMAT    TP1 0009
C AND SETUP OF THESE CARDS SHOULD BE CLEAR FROM THE READ   TP1 0010

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```

C STATEMENTS AT STATEMENT LABEL 2, IN LOOP 7 AND LOOP 10.
C
C     INTEGER VWT(30,30),VFWT(30,60),POST(30)
C
C FOR EXPLANATION OF THESE VECTORS, SEE SUBROUTINE 'LOCATE'.
C
C     READ(5,1) NPRLM
C     1 FORMAT(I3)
C     NP=1
C     2 READ(5,3) NVPT,NFPT
C     3 FORMAT(2I3)
C     WRITE(6,4) NP,NVPT,NFPT
C     4 FORMAT(1H1////14H TEST PROBLEM ,I3//11H NUMBER OF ,
C     119H VARIABLE POINTS IS ,I3,24H;    NUMBER OF FIXED POI,
C     17HN TS IS ,I2)
C     WRITE(6,5)
C     5 FORMAT(20H0WEIGHT MATRIX VWT :/)
C     DO 7 I=1,NVPT
C     READ(5,6) (VWT(I,J),J=1,NVPT)
C     6 FORMAT(25I3)
C     7 WRITE(6,8) (VWT(I,J),J=1,NVPT)
C     8 FORMAT(5X,22(I3,1X)/13X,20(I3,1X)/13X,20(I3,1X))
C     WRITE(6,9)
C     9 FORMAT(21H0WEIGHT MATRIX VFWT :/)
C     DO 10 I=1,NVPT
C     READ(5,6) (VFWT(I,J),J=1,NFPT)
C     10 WRITE(6,8) (VFWT(I,J),J=1,NFPT)
C     CALL LOCATE(NVPT,NFPT,VWT,VFWT,POST)
C     WRITE(6,11) (POST(I),I=1,NVPT)
C     11 FORMAT(38H-OPTIMAL LOCATIONS OF VARIABLE POINTS ,
C     142H COINCIDE WITH THE FOLLOWING FIXED POINTS ://3X,
C     122(I3,1X)/3X,20(I3,1X))
C
C CHECK WHETHER THERE IS MORE TEST PROBLEM
C     NP=NP+1
C     IF(NP.LE.NPRLM) GO TO 2
C     STOP
C     END
C
C *TEST PROGRAM #2* (DATA GENERATED RANDOMLY)
C THIS IS THE SECOND OF THE TWO MAIN PROGRAMS FOR TESTING
C THE SUBROUTINE 'LOCATE'. IT TESTS DATA GENERATED RANDOMLY
C BY A SUBROUTINE CALLED RANDU. IT FIRST READS AN INTEGER
C NPRLM FROM A CARD, USING FORMAT(I3), WHERE NPRLM IS THE
C NUMBER OF PROBLEMS TO BE TESTED IN A SINGLE RUN. FOR EACH
C PROBLEM TESTED, THE USER PROVIDES ONLY A SINGLE CARD TO BE
C READ BY THE FOLLOWING STATEMENTS:
C
C     2 READ(5,3) NVPT,NFPT,NCHECK,MVWT,MVFWT,RANGE
C
C     3 FORMAT(5I3,F6.2)
C
C WHERE
C NVPT - NUMBER OF VARIABLE POINTS.
C NFPT - NUMBER OF FIXED POINTS.
C NCHECK - NUMBER OF SETS OF VARIABLE-POINT LOCATIONS
C GENERATED FOR CHECKING THE OPTIMALITY OF THE
C OBTAINED SOLUTION.
C MVWT - UPPER BOUND ON THE WEIGHTS (F1000) BETWEEN THE
C VARIABLE POINTS.
C MVFWT - UPPER BOUND ON THE WEIGHTS (F1000) BETWEEN THE
C FIXED AND VARIABLE POINTS.
C RANGE - UPPER BOUND ON THE RANGE OF LOCATIONS GENERATED.
C
C USING THESE INPUT DATA, THIS PROGRAM FIRST RANDOMLY GENERATES
C A SET OF FIXED-POINTS LOCATIONS. IT THEN GENERATES RANDOMLY
C NCHECK SETS OF VARIABLE-POINT LOCATIONS AND COMPARES THEM
C WITH THE OPTIMAL LOCATIONS OBTAINED BY 'LOCATE'.
C THIS MAIN PROGRAM INCLUDES THREE OTHER SUBROUTINES:
C RANDU, CHECK AND SUM.
C
C     INTEGER VWT(30,30),VFWT(30,60),POST(30)

```

```

TP1 0011
TP1 0012
TP1 0013
TP1 0014
TP1 0015
TP1 0016
TP1 0017
TP1 0018
TP1 0019
TP1 0020
TP1 0021
TP1 0022
TP1 0023
TP1 0024
TP1 0025
TP1 0026
TP1 0027
TP1 0028
TP1 0029
TP1 0030
TP1 0031
TP1 0032
TP1 0033
TP1 0034
TP1 0035
TP1 0036
TP1 0037
TP1 0038
TP1 0039
TP1 0040
TP1 0041
TP1 0042
TP1 0043
TP1 0044
TP1 0045
TP1 0046
TP1 0047
TP1 0048
C
TP2 0000
TP2 0001
TP2 0002
TP2 0003
TP2 0004
TP2 0005
TP2 0006
TP2 0007
TP2 0008
TP2 0009
TP2 0010
TP2 0011
TP2 0012
TP2 0013
TP2 0014
TP2 0015
TP2 0016
TP2 0017
TP2 0018
TP2 0019
TP2 0020
TP2 0021
TP2 0022
TP2 0023
TP2 0024
TP2 0025
TP2 0026
TP2 0027
TP2 0028
TP2 0029
TP2 0030
TP2 0031
TP2 0032
TP2 0033
TP2 0034

```

C		TP2 0035
C	FOR EXPLANATION OF THESE VECTORS, SEE SUBROUTINE 'LOCATE'	TP2 0036
C		TP2 0037
	READ(5,1) NPRLM	TP2 0038
	1 FORMAT(I3)	TP2 0039
	IY=2*NPRLM+1	TP2 0040
	NP=1	TP2 0041
	2 READ(5,3) NVPT,NFPT,NCHECK,MVWT,MVFWT,RANGE	TP2 0042
	3 FORMAT(5I3,F6.2)	TP2 0043
	WRITE(6,4) NP,NVPT,NFPT	TP2 0044
	4 FORMAT(1H1////14H TEST PROBLEM ,I3//11H NUMBER OF ,	TP2 0045
	119HVARIABLE POINTS IS ,I3,24H; NUMBER OF FIXED POI,	TP2 0046
	17HNPTS IS ,I2)	TP2 0047
	WRITE(6,5) MVWT	TP2 0048
	5 FORMAT(35H0WEIGHT MATRIX VWT (UPPER BOUND IS ,I3,2H):/)	TP2 0049
C		TP2 0050
C	GENERATE MATRIX VWT CONTAINING WEIGHTS BETWEEN VARIABLE	TP2 0051
C	POINTS.	TP2 0052
	DO 9 I=1,NVPT	TP2 0053
	DO 6 J=I,NVPT	TP2 0054
	IX=IY	TP2 0055
	CALL RANDU(IX,IY,YFL)	TP2 0056
	TEMP=FLOAT(MVWT)*YFL	TP2 0057
	VWT(I,J)=INT(TEMP)	TP2 0058
	VWT(J,I)=VWT(I,J)	TP2 0059
	IF (I.EQ.J) VWT(I,J)=0	TP2 0060
	6 CONTINUE	TP2 0061
	WRITE(6,7) (VWT(I,J),J=1,NVPT)	TP2 0062
	7 FORMAT(5X,22(I3,1X)/13X,20(I3,1X)/13X,20(I3,1X))	TP2 0063
C		TP2 0064
C	GENERATE MATRIX VFWT CONTAINING WEIGHTS BETWEEN VARIABLE AND	TP2 0065
C	FIXED POINTS.	TP2 0066
	DO 8 J=1,NFPT	TP2 0067
	IX=IY	TP2 0068
	CALL RANDU(IX,IY,YFL)	TP2 0069
	TEMP=FLOAT(MVFWT)*YFL	TP2 0070
	VFWT(I,J)=INT(TEMP)	TP2 0071
	8 CONTINUE	TP2 0072
	9 CONTINUE	TP2 0073
	WRITE(6,10) MVFWT	TP2 0074
	10 FORMAT(36H0WEIGHT MATRIX VFWT (UPPER BOUND IS ,I3,	TP2 0075
	12H):/)	TP2 0076
	DO 11 I=1,NVPT	TP2 0077
	11 WRITE(6,7) (VFWT(I,J),J=1,NFPT)	TP2 0078
C		TP2 0079
C	FIND AN OPTIMAL SOLUTION BY CALLING 'LOCATE'.	TP2 0080
	CALL LOCATE(NVPT,NFPT,VWT,VFWT,POST)	TP2 0081
	WRITE(6,12) (POST(I),I=1,NVPT)	TP2 0082
	12 FORMAT(38H-OPTIMAL LOCATIONS OF VARIABLE POINTS ,	TP2 0083
	142HCOINCIDE WITH THE FOLLOWING FIXED POINTS ://3X,	TP2 0084
	122(I3,1X)/3X,20(I3,1X))	TP2 0085
C		TP2 0086
C	TO TEST THE OPTIMALTY OF THE SOLUTION OBTAINED BY 'LOCATE',	TP2 0087
C	THE SUBROUTINE CHECK GENERATES NCHECK SETS OF VARIABLE-	TP2 0088
C	POINT LOCATIONS IN THE RANGE (0,RANGE) FOR COMPARISON.	TP2 0089
	CALL CHECK(NVPT,NFPT,VWT,VFWT,POST,NCHECK,RANGE)	TP2 0090
	NP=NP+1	TP2 0091
	IF(NP.LE.NPRLM) GO TO 2	TP2 0092
	STOP	TP2 0093
	END	TP2 0094
	SUBROUTINE RANDU(IX,IY,YFL)	TP2 0095
C		TP2 0096
C	THIS GENERATOR GENERATES A NUMBER RANDOMLY AND UNIFORMLY	TP2 0097
C	IN (0,1). SEE MACLAREN AND MARSALIA, J. ACM, VOL. 12,	TP2 0098
C	PP.83-89, AND IBM SCIENTIFIC SUBROUTINE PACKAGE.	TP2 0099
C		TP2 0100
	IY=IX*65539	TP2 0101
	IF (IY) 1,2,2	TP2 0102
	1 IY=IY+2147483647+1	TP2 0103
	2 YFL=IY	TP2 0104
	YFL=YFL*0.4656613E-9	TP2 0105
	RETURN	TP2 0106
	END	TP2 0107

SUBROUTINE CHECK(NVPT,NFPT,VWT,VFWT,POST,NCHECK,RANGE)	TP2 0108
C	TP2 0109
C THIS SUBROUTINE FIRST RANDOMLY GENERATES A SET OF FIXED-	TP2 0110
C POINT LOCATIONS. IT THEN GENERATES RANDOMLY NCHECK SETS OF	TP2 0111
C VARIABLE-POINT LOCATIONS AND COMPARES THEM WITH THE	TP2 0112
C OPTIMAL LOCATIONS OBTAINED BY THE ALGORITHM.	TP2 0113
C	TP2 0114
INTEGER VWT(30,1),VFWT(30,1),POST(1)	TP2 0115
DIMENSION VPLOC(30),FPLOC(60)	TP2 0116
C	TP2 0117
C VPLOC - VARIABLE-POINT LOCATIONS.	TP2 0118
C FPLOC - FIXED-POINT LOCATIONS.	TP2 0119
C THE MEANINGS OF THE OTHER VARIABLES ARE GIVEN IN THE TEST	TP2 0120
C PROGRAM #2.	TP2 0121
C	TP2 0122
C GENERATE FIXED-POINT LOCATIONS.	TP2 0123
C	TP2 0124
IY=NCHECK*2+1	TP2 0125
FPLOC(1)=0.	TP2 0126
DO 1 I=2,NFPT	TP2 0127
IX=IY	TP2 0128
CALL RANDU(IX,IY,YFL)	TP2 0129
FPLOC(I)=YFL*RANGE	TP2 0130
1 CONTINUE	TP2 0131
N1=NFPT-1	TP2 0132
DO 3 I=2,N1	TP2 0133
I1=I+1	TP2 0134
DO 2 J=I1,NFPT	TP2 0135
IF (FPLOC(I).LE.FPLOC(J)) GO TO 2	TP2 0136
T=FPLOC(I)	TP2 0137
FPLOC(I)=FPLOC(J)	TP2 0138
FPLOC(J)=T	TP2 0139
2 CONTINUE	TP2 0140
3 CONTINUE	TP2 0141
T=FPLOC(NFPT)	TP2 0142
WRITE(6,4) RANGE,T	TP2 0143
4 FORMAT(14H0RANDOM CHECK://3X,13HFIXED-POINT ,	TP2 0144
150H LOCATIONS ARE GENERATED RANDOMLY IN THE RANGE (0,	TP2 0145
14H.00,,F6.2,1H)/3X,29HVARIABLE-POINT LOCATIONS ARE ,	TP2 0146
138HGENERATED RANDOMLY IN THE RANGE (0.00,,F6.2,1H)//	TP2 0147
14X,26HFIXED-POINT LOCATIONS ARE:)	TP2 0148
WRITE(6,5) (FPLOC(I),I=1,NFPT)	TP2 0149
5 FORMAT(7X,12F7.2)	TP2 0150
C	TP2 0151
C GENERATE VARIABLE-POINT LOCATIONS AND EVALUATE THE	TP2 0152
C WEIGHTED SUM OF DISTANCES.	TP2 0153
DO 6 I=1,NVPT	TP2 0154
J=POST(I)	TP2 0155
VPLOC(I)=FPLOC(J)	TP2 0156
6 CONTINUE	TP2 0157
WRITE(6,7)	TP2 0158
7 FORMAT(42H0 OPTIMAL VARIABLE-POINT LOCATIONS ARE:)	TP2 0159
WRITE(6,5) (VPLOC(I),I=1,NVPT)	TP2 0160
S=SUM(NVPT,NFPT,VWT,VFWT,VPLOC,FPLOC)	TP2 0161
WRITE(6,8) S	TP2 0162
8 FORMAT(8X,29HWEIGHTED SUM OF DISTANCES IS:,3X,E11.3)	TP2 0163
DO 11 I=1,NCHECK	TP2 0164
DO 9 J=1,NVPT	TP2 0165
IX=IY	TP2 0166
CALL RANDU(IX,IY,YFL)	TP2 0167
VPLOC(J)=YFL*T	TP2 0168
9 CONTINUE	TP2 0169
WRITE(6,10) I	TP2 0170
10 FORMAT(15H0 CHECK #(,I3,2H):,5X,11HVARIABLE-PO,	TP2 0171
1 18HINT LOCATIONS ARE:)	TP2 0172
WRITE(6,5) (VPLOC(J),J=1,NVPT)	TP2 0173
S=SUM(NVPT,NFPT,VWT,VFWT,VPLOC,FPLOC)	TP2 0174
WRITE(6,8) S	TP2 0175
11 CONTINUE	TP2 0176
RETURN	TP2 0177
END	TP2 0178
FUNCTION SUM(NVPT,NFPT,VWT,VFWT,VPLOC,FPLOC)	TP2 0179
C	TP2 0180
C THIS SUBROUTINE CALCULATES THE WEIGHTED SUM OF DISTANCES	TP2 0181


```

C BETWEEN ALL THE VARIABLE POINTS AND FIXED POINTS. TP2 0182
C THE MEANINGS OF THE INPUT PARAMETERS ARE EXPLAINED IN THE TP2 0183
C SUBROUTINE 'CHECK' AND TEST PROGRAM #2. TP2 0184
C TP2 0185
      INTEGER VWT(30,1),VFWT(30,1) TP2 0186
      DIMENSION VPLOC(1),FPLOC(1) TP2 0187
      SUM=0. TP2 0188
      DO 3 I=1,NVPT TP2 0189
        DO 1 J=1,I TP2 0190
          IF (I.NE.J) SUM=SUM+FLOAT(VWT(I,J)) TP2 0191
1        *ABS(VPLOC(I)-VPLOC(J)) TP2 0192
1      CONTINUE TP2 0193
        DO 2 J=1,NFPT TP2 0194
          SUM=SUM+FLOAT(VFWT(I,J))*ABS(VPLOC(I)-FPLOC(J)) TP2 0195
2      CONTINUE TP2 0196
3      CONTINUE TP2 0197
      RETURN TP2 0198
      END TP2 0199

      SUBROUTINE LOCATE(NVPT,NFPT,VWT,VFWT,POST) LOC 0000
C LOC 0001
C THIS SUBROUTINE SOLVES THE FOLLOWING ONE-DIMENSIONAL LOC 0002
C MULTIFACILITY LOCATION PROBLEM USING AN ALGORITHM DESIGNED LOC 0003
C BY PICARD, RATLIFF (SEE OPERATIONS RESEARCH 26 (1978), LOC 0004
C PP.422-433) AND CHEUNG. LOC 0005
C " GIVEN NFPT FIXED POINTS ON A LINE, DETERMINE THE LOC 0006
C LOCATIONS OF NVPT VARIABLE POINTS SUCH THAT A WEIGHTED LOC 0007
C SUM OF THE DISTANCES BETWEEN THE POINTS IS MINIMUM." LOC 0008
C OUR ALGORITHM TRANSFORMS THE LOCATION PROBLEM TO A SEQUENCE LOC 0009
C OF MINIMUM-CUT PROBLEMS. A FORTRAN CODE (THE SUBROUTINE LOC 0010
C NETFLO IN "COMBINATORIAL ALGORITHMS", WRITTEN BY NIJENHUIS LOC 0011
C AND WILF) OF DINIC'S ALGORITHM IS USED TO FIND THE MINIMUM LOC 0012
C CUTS. LOC 0013
C TO APPLY THE SUBROUTINE 'LOCATE', THE FIXED POINTS SHOULD LOC 0014
C BE SUBSCRIPTED ON THE LINE IN INCREASING ORDER OF THEIR LOC 0015
C LOCATIONS. THEIR EXACT LOCATIONS ARE NOT REQUIRED. LOC 0016
C LOC 0017
      INTEGER COL,SOURCE,SUM LOC 0018
      INTEGER POST(1),CVPT(32),CUT(32),ENDPT(4,930) LOC 0019
      INTEGER VWT(30,1),VFWT(30,1) LOC 0020
C LOC 0021
C ON INPUT LOC 0022
C NVPT NUMBER OF VARIABLE POINTS. LOC 0023
C NFPT NUMBER OF FIXED POINTS. LOC 0024
C VWT A SQUARE MATRIX WHOSE UPPER TRIANGULAR PART LOC 0025
C CONTAINS THE WEIGHTS BETWEEN THE VARIABLE POINTS. LOC 0026
C THE WEIGHTS MUST BE NON-NEGATIVE INTEGERS. LOC 0027
C VFWT WEIGHTS BETWEEN THE VARIABLE AND FIXED POINTS. LOC 0028
C THE WEIGHTS MUST BE NON-NEGATIVE INTEGERS. LOC 0029
C ON OUTPUT LOC 0030
C POST A VECTOR INDICATING THE OPTIMAL LOCATIONS OF THE LOC 0031
C VARIABLE POINTS. VARIABLE POINT I WILL COINCIDE LOC 0032
C WITH THE FIXED POINT WHOSE POSITION IS POST(I). LOC 0033
C FOR WORKING STORAGE LOC 0034
C (IN THE FOLLOWING, MNVP MEANS 'THE MAXIMUM NUMBER OF LOC 0035
C VARIABLE POINTS'.) LOC 0036
C CVPT A VECTOR OF LENGTH MNVP+2. IT REPRESENTS THE LOC 0037
C VARIABLE POINTS WHICH PARTICIPATE IN THE CURRENT LOC 0038
C ITERATION. THE VARIABLE POINT CVPT(I) IS THE LOC 0039
C NODE I OF THE CURRENT NETWORK. LOC 0040
C ENDPT A MATRIX OF DIMENSION 4 X E, WHERE E=MNVP* LOC 0041
C *(MNVP+1). EACH ARC CORRESPONDS TO A COLUMN OF LOC 0042
C ENDPT, WHOSE 4 ROWS DENOTE THE INITIAL VERTEX, LOC 0043
C THE TERMINAL VERTEX, THE CAPACITY AND MAXIMUM LOC 0044
C FLOW OF THAT ARC RESPECTIVELY. LOC 0045
C CUT A VECTOR OF LENGTH MNVP+2. FOR A MIN-CUT, LOC 0046
C CUT(I)=1 MEANS NODE I IS ASSOCIATED WITH THE LOC 0047
C SOURCE, AND CUT(I)=0 MEANS NODE I IS ASSOCIATED LOC 0048
C WITH THE SINK. LOC 0049
C LOC 0050
      IF (NVPT.EQ.1) GO TO 13 LOC 0051
C LOC 0052
C INITIALIZE ENDPT FOR THE FIRST ITERATION. LOC 0053
      SOURCE=NVPT+1 LOC 0054
      COL=0 LOC 0055

```


Q=VERT(P, I)	NET 0058
IF(CAP(P, Q).EQ.0) GO TO 12	NET 0059
IF(Q.EQ.SINK) LBLSNK=-LABEL	NET 0060
IF(CUT(Q)-LABEL) 10, 11, 12	NET 0061
10 CUT(Q)=LABEL	NET 0062
WR=WR+1	NET 0063
AUX(WR)=Q	NET 0064
11 I=I+1	NET 0065
GO TO 9	NET 0066
12 VERT(P, I)=VERT(P, M)	NET 0067
VERT(P, M)=Q	NET 0068
M=M-1	NET 0069
GO TO 9	NET 0070
13 CUT(P)=M	NET 0071
RD=RD+1	NET 0072
IF(RD.GT.WR) GO TO 24	NET 0073
P=AUX(RD)	NET 0074
IF(CUT(P)+LBLSNK.EQ.0) GO TO 14	NET 0075
LABEL=CUT(P)-1	NET 0076
GO TO 8	NET 0077
C	NET 0078
C CONSTRUCTION OF PATH FROM SOURCE TO SINK	NET 0079
14 Q=SOURCE	NET 0080
K=0	NET 0081
15 K=K+1	NET 0082
AUX(K)=Q	NET 0083
IF(K.GT.LBLSNK) GO TO 19	NET 0084
16 P=AUX(K)	NET 0085
17 M=CUT(P)	NET 0086
IF(M.EQ.0) GO TO 18	NET 0087
Q=VERT(P, M)	NET 0088
GO TO 15	NET 0089
18 K=K-1	NET 0090
IF(K.EQ.0) GO TO 6	NET 0091
P=AUX(K)	NET 0092
CUT(P)=CUT(P)-1	NET 0093
GO TO 17	NET 0094
19 IF(Q.NE.SINK) GO TO 18	NET 0095
DELTA=CAP(P, Q)	NET 0096
DO 20 I=2, K	NET 0097
I1=AUX(I-1)	NET 0098
I2=AUX(I)	NET 0099
DELTA=MIN0(DELTA, CAP(I1, I2))	NET 0100
20 CONTINUE	NET 0101
21 K=K-1	NET 0102
IF(K.EQ.0) GO TO 23	NET 0103
P=AUX(K)	NET 0104
C=CAP(P, Q)-DELTA	NET 0105
IF(C.GT.0) GO TO 22	NET 0106
CUT(P)=CUT(P)-1	NET 0107
KO=K	NET 0108
22 CAP(P, Q)=C	NET 0109
CAP(Q, P)=CAP(Q, P)+DELTA	NET 0110
Q=P	NET 0111
GO TO 21	NET 0112
23 FLOVAL=FLOVAL+DELTA	NET 0113
K=KO	NET 0114
GO TO 16	NET 0115
C	NET 0116
C EXIT PROCEDURE.	NET 0117
24 DO 25 I=1, N	NET 0118
CUT(I)=MIN0(1, CUT(I)-NMIN)	NET 0119
25 CONTINUE	NET 0120
DO 26 I=1, E	NET 0121
I1=ENDPT(1, I)	NET 0122
I2=ENDPT(2, I)	NET 0123
ENDPT(4, I)=ENDPT(3, I)-CAP(I1, I2)	NET 0124
26 CONTINUE	NET 0125
RETURN	NET 0126
END	NET 0127

18

3 4
0 3 2
3 0 1

```

2 1 0
3 2 3 1
1 2 1 2
9 1 2 2
3 4
0 1 3
1 0 3
3 3 0
9 1 1 2
2 1 3 2
2 2 4 1
5 20
0 21 14 9 7
21 0 34 24 32
14 34 0 14 36
9 24 14 0 8
7 32 36 8 0
50 49 81 46 52 80 83 48 8 33 99 0 50 76 38 40 70 32 30 92
71 63 77 31 62 5 2 85 13 37 19 64 17 59 49 75 52 44 88 94
73 32 27 72 2 7 66 79 1 19 22 76 50 68 56 87 83 69 58 36
97 24 64 49 52 1 74 8 88 84 53 84 46 74 92 45 29 19 54 71
26 51 46 49 93 34 46 35 19 28 21 15 24 31 41 89 22 92 97 99
5 20
0 21 14 9 7
21 0 34 24 32
14 34 0 14 36
9 24 14 0 8
7 32 36 8 0
33 0 50 48 46 38 32 92 80 99 30 8 40 70 52 50 49 83 76 81
37 64 71 85 31 49 44 94 5 19 88 13 75 52 62 17 63 2 59 77
19 76 73 79 72 56 69 36 7 22 58 1 87 83 2 50 32 66 68 27
84 84 97 8 49 92 19 71 1 53 54 88 45 29 52 46 24 74 74 64
28 15 26 35 49 41 92 99 34 21 97 19 89 22 93 24 51 46 31 46
5 3
0 2 2 2 2
2 0 20 1 0
2 20 0 0 0
2 1 0 0 40
2 0 0 40 0
10 0 0
4 1 4
4 1 4
4 5 4
4 5 4
3 5
0 1 1
1 0 1
1 1 0
1 1 6 1 6
4 1 1 1 1
1 1 1 1 1
3 5
0 1 0
1 0 1
0 1 0
3 2 2 0 0
2 2 2 4 4
4 0 2 6 4
3 5
0 1 0
1 0 1
0 1 0
2 3 0 2 0
2 2 4 2 4
2 4 6 0 4
1 15
1
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
1 15
1
1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
1 20
1
4 2 15 9 0 45 3 7 5 0 0 1 50 4 0 12 32 4 22 7
5 15
0 1 1 1 1

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1 0 1 1 1
1 1 0 1 1
1 1 1 0 1
1 1 1 1 0
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
5 15
0 1 1 1 1
1 0 1 1 1
1 1 0 1 1
1 1 1 0 1
1 1 1 1 0
1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
6 20
0 5 2 3 12 8
5 0 6 40 8 6
2 6 0 1 18 7
3 40 1 0 9 13
12 8 18 9 0 4
8 6 7 13 4 0
4 2 15 9 0 45 3 7 5 0 0 1 50 4 0 12 32 4 22 7
4 2 15 9 0 45 3 7 5 0 0 1 50 4 0 12 32 4 22 7
4 2 15 9 0 45 3 7 5 0 0 1 50 4 0 12 32 4 22 7
4 2 15 9 0 45 3 7 5 0 0 1 50 4 0 12 32 4 22 7
4 2 15 9 0 45 3 7 5 0 0 1 50 4 0 12 32 4 22 7
4 2 15 9 0 45 3 7 5 0 0 1 50 4 0 12 32 4 22 7
6 20
0 5 2 3 12 8
5 0 6 40 8 6
2 6 0 1 18 7
3 40 1 0 9 13
12 8 18 9 0 4
8 6 7 13 4 0
0 0 0 1 2 0 4 0 0 0 9 5 12 0 2 8 0 2 80 99
0 0 1 0 7 8 2 4 9 2 0 10 4 6 8 6 0 10 24 73
0 2 1 2 0 3 4 5 0 1 6 9 1 0 5 2 0 4 0 153
9 8 0 0 0 0 0 4 3 0 5 3 0 1 1 1 0 6 49 88
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 22
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
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0 5 2 3 12 8
5 0 6 40 8 6
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3 40 1 0 9 13
12 8 18 9 0 4
8 6 7 13 4 0
0 0 1 0 7 8 2 4 9 2 98 10 4 6 8 6 0 1 1 2
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
0 0 0 1 2 0 4 0 0 111 9 5 12 0 2 8 0 2 0 0
0 2 1 2 0 3 4 5 0 57 6 9 1 0 5 2 0 4 0 8
9 8 0 0 0 0 0 4 3 10 25 3 0 1 1 1 0 6 0 0
24 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
6 20
0 1 1 1 1 1
1 0 1 1 1 1
1 1 0 1 1 1
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1 1 1 1 0 1
1 1 1 1 1 0
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147 0 0 1 2 0 4 0 0 0 9 5 12 0 2 8 0 2 0 0
280 0 1 0 7 8 2 4 9 2 0 10 4 6 8 6 0 2 3 0
230 22 0 0 0 0 0 4 3 0 5 3 0 1 1 1 0 6 1 0
30 20 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 4
4 99 1 2 0 3 4 5 0 1 6 9 1 0 5 2 0 4 0 0
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
6 20
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1 1 0 1 1 1

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1 1 1 0 1 1
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280 0 1 0 7 8 2 4 9 2 0 10 4 6 8 6 0 2 3 0
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
4 99 1 2 0 3 4 5 0 1 6 9 1 0 5 2 0 4 0 0
30 20 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 4
147 0 0 1 2 0 4 0 0 0 9 5 12 0 2 8 0 2 0 0
230 22 0 0 0 0 0 4 3 0 5 3 0 1 1 1 0 6 1 0
6 20
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4 2 15 9 0 45 3 7 5 0 0 1 50 4 0 12 32 4 22 7
0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 1 1 1
30 20 1 1 1 1 22 15 0 3 1 1 1 1 1 1 21 1 1 4
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
0 0 0 1 2 0 4 0 0 0 9 5 12 0 2 8 0 2 0 24
4 99 1 2 0 3 4 5 0 1 6 9 1 0 5 2 0 4 0 0

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7
2 5 10 50500100.00
3 5 10 30400200.00
3 4 10 50500150.00
10 20 20 50500200.00
15 20 25 50500300.00
20 30 30 25400250.00
30 50 50 30500500.00

```


ALGORITHM 559

The Stationary Point of a Quadratic Function Subject to Linear Constraints [E4]

J. T. BETTS
The Aerospace Corporation

Key Words and Phrases: quadratic programming, orthogonal decomposition
CR Categories: 5.14, 5.15, 5.4, 5.40, 5.41
Language: Fortran

DESCRIPTION

This algorithm implements the method developed in [1].

REFERENCE

1. BETTS, J.T. A compact algorithm for computing the stationary point of a quadratic function subject to linear constraints. *ACM Trans. Math. Softw.* 6, 3 (Sept. 1980), 391-397.

ALGORITHM

```

C      PROGRAM DRIVER(INPUT,OUTPUT,TAPE5=INPUT,TAPE6=OUTPUT)
C      DIMENSION A(4,4), B(4), C(4,4), D(4), G(4), H(4), U(4), IP(4),
C      * X(4)
C      DIMENSION DJNORM(1)
C      DIMENSION DMCX(4), AA(4,4), BB(4), CC(4,4), DD(4)
C      DIMENSION NN(12), MM(12), KEYKEY(12)
C      DATA NN /0,1,1,3,3,3,3,3,4,4,4,4/
C      DATA MM /0,0,1,0,1,2,3,4,0,3,0,2/
C      DATA KEYKEY /1,1,1,1,1,1,1,1,2,2,3,3/
C      DATA MAXRA, MAXRC, MAXCAS /4,4,12/
C      DATA TAU /1.E-12/
C
C      THE PURPOSE OF THIS DRIVER IS TO DEMONSTRATE THE USE OF
C      SUBROUTINE HSQP
C
C      TEST CASES ARE SPECIFIED BY THE CONTENTS OF THE
C      ARRAYS NN(), MM(), AND KEYKEY().
C      KEY = 1 SETS A AND C NONSINGULAR.
C      = 2 SETS RANK OF A = MIN( N, 2)
C      AND RANK OF C = MIN( M, N, 2)
C      = 3 SETS C NONSINGULAR AND ALL ELTS OF A = 0.
C
C      SET TAU TO ABOUT 100 TIMES THE MACHINE PRECISION.
C
C      DO 230 KASE=1,MAXCAS
C      N = NN(KASE)
C      M = MM(KASE)
C      KEY = KEYKEY(KASE)

```

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Author's address: The Aerospace Corporation, P.O. Box 92957, Los Angeles, CA 90009.

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```

        WRITE (6,1001) KASE, N, M, KEY, TAU
C
C         DEFINE OBJECTIVE FUNCTION
C
        IF (N.LE.0) GO TO 160
        DO 60 I=1,N
        B(I) = FLOAT(I) - .05 * FLOAT(I)**2
        BB(I) = B(I)
        DO 50 J=I,N
        GO TO (10, 20, 30), KEY
100   A(I,J) = 1./(FLOAT(J-I)+.3+.01*FLOAT(I))
        GO TO 40
200   A(I,J) = FLOAT(I+J)
        GO TO 40
300   A(I,J) = 0.
400   CONTINUE
        A(J,I) = A(I,J)
        AA(I,J) = A(I,J)
        AA(J,I) = A(J,I)
500   CONTINUE
600   CONTINUE
C
C         DEFINE CONSTRAINTS
C
        IF (M.LE.0) GO TO 120
        DO 110 I=1,M
        D(I) = FLOAT(I+1) - .05 * FLOAT(I+1)
        DD(I) = D(I)
        DO 100 J=1,N
        C(I,J) = FLOAT(I-J+4)
        GO TO (70, 80, 70), KEY
700   C(I,J) = 1./(FLOAT(J-I)+.2+.01*FLOAT(I+J))
        GO TO 90
800   C(I,J) = FLOAT(I-J+4)
900   CONTINUE
        CC(I,J) = C(I,J)
1000  CONTINUE
1100  CONTINUE
1200  CONTINUE
        WRITE (6,1002)
        DO 130 I=1,N
        WRITE (6,1008) (A(I,J),J=1,N), B(I)
1300  CONTINUE
        IF (M.LE.0) GO TO 150
        WRITE (6,1003)
        DO 140 I=1,M
        WRITE (6,1008) (C(I,J),J=1,N), D(I)
1400  CONTINUE
1500  CONTINUE
1600  CONTINUE
C
C         CALL HSQP TO SOLVE PROBLEM
C
        CALL HSQP(A, B, C, D, M, N, TAU, G, H, U, IP, MAXRA, MAXRC,
* DJNORM, X, KRANK)
C
C         WRITE SOLUTION VECTOR
C
        WRITE (6,1007) DJNORM(1), KRANK
        IF (DJNORM(1).LT.0. .OR. N.LE.0) GO TO 220
        WRITE (6,1006) (X(I),I=1,N)
C
C         CHECK CONSTRAINTS.
C
        IF (M.LE.0) GO TO 190
        DO 180 I=1,M
        DMCX(I) = DD(I)
        DO 170 J=1,N
        DMCX(I) = DMCX(I) - CC(I,J)*X(J)
1700  CONTINUE
1800  CONTINUE
        WRITE (6,1004) (DMCX(I),I=1,M)
1900  CONTINUE
C
C         EVALUATE OBJECTIVE FUNCTION.
C

```

```

      VALUE = 0.
      DO 210 J=1,N
      SUM = 0.
      DO 200 I=1,N
      SUM = SUM + X(I)*AA(I,J)
200 CONTINUE
      VALUE = VALUE + (0.5*SUM+BB(J))*X(J)
210 CONTINUE
      WRITE (6,1005) VALUE
220 CONTINUE
230 CONTINUE
      STOP
1001 FORMAT (51H0*****
*/31H CASE N M KEY TAU/1X,4I5,E10.3)
1002 FORMAT (26H0 AUGMENTED MATRIX (A:B) =)
1003 FORMAT (26H0 AUGMENTED MATRIX (C:D) =)
1004 FORMAT (33H0 CONSTRAINT RESIDUALS D - C*X =/(1X, 8E14.3))
1005 FORMAT (32H0 VALUE OF OBJECTIVE FUNCTION = , E20.6)
1006 FORMAT (16H0 SOLUTION VECTOR/(1X, 8E14.6))
1007 FORMAT (24H0 PROJECTED GRADIENT NORM, 5X, E20.8/ 13H0 RANK OF PROJ,
* 20HECTED HESSIAN MATRIX, 5X, 110)
1008 FORMAT (1X, 8E14.6)
      END

```

```

      SUBROUTINE HSQP(A,B,C,D,M,N,TAU,G,H,U,IP,MAXRA,MAXRC,DJNORM,X,
      $ KRANK)

```

C
C

```

      DIMENSION B(1),D(1),G(1),H(1),U(1),IP(1),DJNORM(1),X(1)
      DIMENSION A(MAXRA,1),C(MAXRC,1)

```

C
C

PROGRAMMER AND DATE: J.T.BETTS, JAN. 1978.

C
C

PURPOSE: GIVEN AN M X N MATRIX C (OF RANK M), AN M VECTOR D,
AN N X N SYMMETRIC MATRIX A, AND AN N VECTOR B, FIND THE
STATIONARY POINT X OF THE QUADRATIC

C
C

$$J = .5*(X**T)*A*X + (B**T)*X$$

C
C

SUBJECT TO THE CONSTRAINTS

C
C

$$C*X = D,$$

C
C

IF A STATIONARY POINT DOES NOT EXIST THE ALGORITHM WILL FIND
A POINT WHICH SATISFIES THE CONSTRAINTS AND MINIMIZES THE
NORM OF THE GRADIENT OF J PROJECTED ON THE CONSTRAINT SURFACE.

C
C

ALGORITHM; ORTHOGONAL DECOMPOSITION OF C MATRIX USING
HOUSEHOLDER TRANSFORMATIONS, FOLLOWED BY APPLICATION OF THE
OPTIMALITY CONDITIONS IN THE REDUCED VARIABLES.

C
C

INPUT:

C
C

A	N X N SYMMETRIC HESSIAN MATRIX
B	N DIMENSIONAL GRADIENT VECTOR
C	M X N JACOBIAN MATRIX (RANK M)
D	M DIMENSIONAL CONSTRAINT VECTOR
M	THE NUMBER OF CONSTRAINTS
N	THE NUMBER OF VARIABLES
TAU	PSEUDORANK TEST PARAMETER. FOR A MACHINE WITH K SIGNIFICANT FIGURES AN APPROPRIATE VALUE IS TAU = 1.E-(K-2).
G	AUXILLIARY STORAGE (LENGTH M)
H	AUXILLIARY STORAGE (LENGTH N-M)
U	AUXILLIARY STORAGE (LENGTH N-M)
IP	AUXILLIARY STORAGE (LENGTH N-M)
MAXRA	MAXIMUM ROW DIMENSION OF A (MAXRA X N)
MAXRC	MAXIMUM ROW DIMENSION OF C (MAXRC X M)

C
C

OUTPUT:

C
C

DJNORM IF DJNORM .GE. 0. IT IS THE NORM OF THE PROJECTED

```

C          GRADIENT OF THE CONSTRAINED QUADRATIC FORM.
C          IF DJNORM = 0. X IS A STATIONARY POINT OF THE
C          CONSTRAINED QUADRATIC FORM.
C          DJNORM = -1. MEANS INPUT ERRORS.      NO X COMPUTED.
C          DJNORM = -2. MEANS RANK(C) .LT. M.    NO X COMPUTED.
C          X          COMPUTED SOLUTION VECTOR. IT IS A STATIONARY POINT
C          IF DJNORM = 0.
C          KRANK      PSEUDORANK OF PROJECTED HESSIAN MATRIX (K2**T)*A*K2.
C          IF KRANK = N-M, X IS THE UNIQUE SOLUTION.
C          IF KRANK .LT. N-M, X IS THE MINIMUM LENGTH SOLUTION.
C
C          NOTE:      THE INPUT VALUES OF A,B,C, AND D ARE DESTROYED.
C
C -----
C
C          INITIALIZATION
C
C          KRANK = 0
C          MP1 = M + 1
C          NMM = N - M
C          DJNORM(1) = -1.
C
C          CHECK FOR INPUT ERRORS
C
C          IF(N.LE.0.OR.N.GT.MAXRA.OR.M.GT.MAXRC.OR.M.GT.N
C          $ .OR. M .LT. 0) RETURN
C
C          COMPUTE TOLERANCE, ATOL, FOR PSEUDORANK OF A.
C
C          TEMP = 0.
C          DO 20 J = 1,N
C          SUMSQ = 0.
C          DO 10 I = 1,N
C          SUMSQ = SUMSQ + A(I,J)**2
10    CONTINUE
C          TEMP = AMAX1(TEMP,SUMSQ)
20    CONTINUE
C          ATOL = TAU*SQRT(TEMP)
C
C          IF THE PROBLEM IS UNCONSTRAINED GO TO STEP 6
C
C          IF(M.EQ.0) GO TO 140
C
C -----
C
C          STEP. 1.  COMPUTE ORTHOGONAL MATRIX K.  TRIANGULARIZE C.
C
C          DO 30 I = 1,M
C          CALL H12(1,I,I+1,N,C(I,1),MAXRC,G(I),C(I+1,1),MAXRC,1,M-I)
30    CONTINUE
C
C -----
C
C          STEP 2.  COMPUTE YIHAT BY SOLVING THE LOWER TRIANGULAR
C          SYSTEM C*Y1 = D.  STORE IN X.
C
C          TEMP = 0.
C          DO 50 J=1,M
C          DO 40 I=J,M
C          TEMP=AMAX1(TEMP,ABS(C(I,J)))
40    CONTINUE
50    CONTINUE
C          CTOL = TAU*TEMP
C          DO 90 I = 1,M
C          IM1 = I - 1
C          X(I) = D(I)
C          IF(I .EQ. 1) GO TO 70
C          DO 60 J = 1,IM1
C          X(I) = X(I) - C(I,J)*X(J)
60    CONTINUE
70    CONTINUE

```

```

      IF (ABS(C(I,I)) .GT. CTOL) GO TO 80
      DJNORM(1) = -2.
      RETURN
80  CONTINUE
      X(I) = X(I)/C(I,I)
90  CONTINUE
C
C      WHEN THERE ARE NO DEGREES OF FREEDOM GO TO STEP 8
C
      IF (M .LT. N) GO TO 100
      DJNORM(1) = 0.
      GO TO 190
100 CONTINUE
C
C -----
C
C      STEP 3. COMPUTE ATILDA = (K**T)*A
C
      DO 110 I = 1,M
      CALL H12(2,I,I+1,N,C(I,1),MAXRC,G(I),A,1,MAXRA,N)
110 CONTINUE
C
C -----
C
C      STEP 4. FORM THE LAST N-M ROWS OF AHAT = ATILDA*K; I.E.
C      COMPUTE A21HAT = (K2**T)*A*K1 AND A22HAT = (K2**T)*A*K2
C
      DO 120 I = 1,M
      CALL H12(2,I,I+1,N,C(I,1),MAXRC,G(I),A(MP1,1),MAXRA,1,NMM)
120 CONTINUE
C
C -----
C
C      STEP 5. COMPUTE BTILDA = (K**T)*B
C
      DO 130 I = 1,M
      CALL H12(2,I,I+1,N,C(I,1),MAXRC,G(I),B,1,1,1)
130 CONTINUE
C
C -----
C
C      STEP 6. COMPUTE B2HAT = -B2TILDA - A21HAT*Y1HAT
C
140 CONTINUE
      DO 170 I = MP1,N
      B(I) = -B(I)
      IF (M .EQ. 0) GO TO 160
      DO 150 J = 1,M
      B(I) = B(I) - A(I,J)*X(J)
150 CONTINUE
160 CONTINUE
170 CONTINUE
C
C -----
C
C      STEP 7. SOLVE A22HAT*Y2 = B2HAT FOR Y2 USING HFTI
C
      CALL HFTI(A(MP1,MP1),MAXRA,NMM,NMM,B(MP1),1,1,ATOL,KRANK,
$      DJNORM,H,U,IP)
C
      DO 180 I = MP1,N
      X(I) = B(I)
180 CONTINUE
C
C      IF THE PROBLEM IS UNCONSTRAINED, RETURN.
C
      IF (M.EQ.0) RETURN
C
C -----

```

```

C
C          STEP 8.  COMPUTE X = K*Y
C
190 CONTINUE
   DO 200 K = 1,M
   I = MP1 - K
   CALL H12(2,I,I+1,N,C(I,1),MAXRC,G(I),X,1,1,1)
200 CONTINUE
C
   RETURN
   END

SUBROUTINE H12(MODE,LPIVOT,L1,M,U,IUE,UP,C,ICE,ICV,NCV)
C
C          ALGORITHM H12:  C.L. LAWSON AND R.J. HANSON, "SOLVING LEAST
C          SQUARES PROBLEMS", PRENTICE-HALL,1974. APPENDIX C,P. 308.
C
C          PURPOSE:  CONSTRUCTION AND/OR APPLICATION OF A SINGLE
C          HOUSEHOLDER TRANSFORMATION ...  $Q = I + U*(U**T)/B$ 
C
C          MODE = 1 OR 2 TO SELECT ALGORITHM H1 OR H2.
C          LPIVOT IS THE INDEX OF THE PIVOT ELEMENT
C          L1,M   IF L1.LE.M THE TRANSFORMATION WILL BE CONSTRUCTED TO
C                ZERO ELEMENTS INDEXED FROM L1 THROUGH M.  IF L1.GT.M
C                THE SUBROUTINE DOES AN IDENTITY TRANSFORMATION.
C          U(),IUE,UP  ON ENTRY TO H1 U() CONTAINS THE PIVOT VECTOR.
C                IUE IS THE STORAGE INCREMENT BETWEEN ELEMENTS.  ON EXIT
C                FROM H1 U() AND UP CONTAIN QUANTITIES DEFINING THE
C                VECTOR U OF THE HOUSEHOLDER TRANSFORMATION.  ON ENTRY TO
C                H2 U() AND UP SHOULD CONTAIN QUANTITIES PREVIOUSLY
C                COMPUTED BY H1.  THESE WILL NOT BE MODIFIED BY H2.
C          C()      ON ENTRY TO H1 OR H2 C() CONTAINS A MATRIX WHICH WILL
C                BE REGARDED AS A SET OF VECTORS TO WHICH THE
C                HOUSEHOLDER TRANSFORMATION IS TO BE APPLIED.  ON EXIT
C                C() CONTAINS THE SET OF TRANSFORMED VECTORS.
C          ICE     STORAGE INCREMENT BETWEEN ELEMENTS OF VECTORS IN C()
C          ICV     STORAGE INCREMENT BETWEEN VECTORS IN C().
C          NCV     NUMBER OF VECTORS IN C() TO BE TRANSFORMED.  IF NCV.LE.0
C                NO OPERATIONS WILL BE DONE ON C().
C
DIMENSION U(IUE,M),C(1)
DOUBLE PRECISION SM,B
C
IF(0.GE.LPIVOT.OR.LPIVOT.GE.L1.OR.L1.GT.M) RETURN
CL = ABS(U(1,LPIVOT))
IF(MODE.EQ.2) GO TO 60
C
C          CONSTRUCT THE TRANSFORMATION.
C
DO 10 J = L1,M
10 CL = AMAX1(ABS(U(1,J)),CL)
IF(CL) 130,130,20
20 CLINV = 1./CL
SM = (DBLE(U(1,LPIVOT))*CLINV)**2
DO 30 J = L1,M
30 SM = SM + (DBLE(U(1,J))*CLINV)**2
C
C          CONVERT DBLE. PREC. SM TO SNGL. PREC. SMI
C
SMI = SM
CL = CL*SQRT(SMI)
IF(U(1,LPIVOT)) 50,50,40
40 CL = -CL
50 UP = U(1,LPIVOT) - CL
U(1,LPIVOT) = CL
GO TO 70
C
C          APPLY THE TRANSFORMATION  $I + U*(U**T)/B$  TO C
C
60 IF(CL) 130,130,70
70 IF(NCV.LE.0) RETURN
B = DBLE(UP)*U(1,LPIVOT)

```

```

C
C          B MUST BE NONPOSITIVE HERE.  IF B=0 RETURN
C
      IF(B) 80,130,130
80    B = 1.D0/B
      I2 = 1 - ICV + ICE*(LPIVOT-1)
      INCR = ICE*(L1-LPIVOT)
      DO 120 J = 1,NCV
      I2 = I2 + ICV
      I3 = I2 + INCR
      I4 = I3
      SM = C(I2)*DBLE(UP)
      DO 90 I = L1,M
      SM = SM + C(I3)*DBLE(U(1,I))
90    I3 = I3 + ICE
      IF(SM) 100,120,100
100   SM = SM*B
      C(I2) = C(I2) + SM*DBLE(UP)
      DO 110 I = L1,M
      C(I4) = C(I4) + SM*DBLE(U(1,I))
110   I4 = I4 + ICE
120   CONTINUE
130   RETURN
      END

      SUBROUTINE HFTI(A,MDA,M,N,B,MDB,NB,TAU,KRANK,RNORM,H,G,IP)
C
C
C          PURPOSE:  SOLVE THE MATRIX LINEAR LEAST SQUARE PROBLEM
C                   MIN NORM(A*X - B)
C                   WHERE A IS MXN,  B IS MXNB,  X IS NXNB.
C
C          REF.  C.L. LAWSON AND R.J. HANSON, "SOLVING LEAST SQUARES PROBLEMS
C                PRENTICE-HALL,1974.  ALGORITHM HFTI, APPENDIX C.,P.290.
C
C
C          DIMENSION A(MDA,N),B(MDB,NB),H(N),G(N),IP(N),RNORM(NB)
C          DOUBLE PRECISION SM,DZERO
C
C          SZERO = 0.
C          DZERO = 0.D0
C          FACTOR = .001
C
C          K = 0
C          LDIAG = MIN0(M,N)
C          IF(LDIAG.LE.0) GO TO 270
C
C          DO 80 J = 1,LDIAG
C          IF(J.EQ.1) GO TO 200
C
C          UPDATE SQUARED COLUMN LENGTHS AND FIND LMAX
C
C          LMAX = J
C          DO 100 L = J,N
C          H(L) = H(L) - A(J-1,L)**2
C          IF(H(L).GT.H(LMAX)) LMAX = L
100   CONTINUE
C          IF(FACTOR*H(LMAX)) 20,20,50
C
C          COMPUTE SQUARED COLUMN LENGTHS AND FIND LMAX
C
C          LMAX = J
C          DO 40 L = J,N
C          H(L) = 0.
C          DO 30 I = J,M
C          H(L) = H(L) + A(I,L)**2
C          IF(H(L).GT.H(LMAX)) LMAX = L
40    CONTINUE
C          HMAX = H(LMAX)
C
C          LMAX HAS BEEN DETERMINED.  DO COLUMN INTERCHANGES IF NEEDED.
C
C          CONTINUE
C          IP(J) = LMAX
C          IF(IP(J).EQ.J) GO TO 70

```

```

      DO 60 I = 1,M
      TMP = A(I,J)
      A(I,J) = A(I,LMAX)
60    A(I,LMAX) = TMP
      H(LMAX) = H(J)
C
C      COMPUTE THE J-TH TRANSFORMATION AND APPLY IT TO A AND B.
C
70    CALL H12(1,J,J+1,M,A(1,J),1,H(J),A(1,J+1),1,MDA,N-J)
80    CALL H12(2,J,J+1,M,A(1,J),1,H(J),B,1,MDB,NB)
C
C      DETERMINE THE PSEUDORANK,K, USING THE TOLERANCE, TAU
C
      DO 90 J = 1,LDIAG
      IF(ABS(A(J,J)).LE.TAU) GO TO 100
90    CONTINUE
      K = LDIAG
      GO TO 110
100   K = J - 1
110   KP1 = K + 1
C
C      COMPUTE THE NORMS OF THE RESIDUAL VECTORS.
C
      IF(NB.LE.0) GO TO 140
      DO 130 JB = 1,NB
      TMP = SZERO
      IF(KP1.GT.M) GO TO 130
      DO 120 I = KP1,M
120   TMP = TMP + B(I,JB)**2
130   RNORM(JB) = SQRT(TMP)
140   CONTINUE
C
C      SPECIAL FOR PSEUDORANK = 0.
C
      IF(K.GT.0) GO TO 160
      IF(NB.LE.0) GO TO 270
      DO 150 JB = 1,NB
      DO 150 I = 1,N
150   B(I,JB) = SZERO
      GO TO 270
C
C      IF THE PSEUDORANK IS LESS THAN N COMPUTE HOUSEHOLDER
C      DECOMPOSITION OF FIRST K ROWS.
C
160   IF(K.EQ.N) GO TO 180
      DO 170 II = 1,K
      I = KP1 - II
170   CALL H12(1,I,KP1,N,A(I,1),MDA,G(I),A,MDA,1,I-1)
180   CONTINUE
C
C
      IF(NB.LE.0) GO TO 270
      DO 260 JB = 1,NB
C
C      SOLVE THE K BY K TRIANGULAR SYSTEM
C
      DO 210 L = 1,K
      SM = DZERO
      I = KP1 - L
      IF(I.EQ.K) GO TO 200
      IP1 = I + 1
      DO 190 J = IP1,K
190   SM = SM + DBLE(A(I,J))*DBLE(B(J,JB))
200   SM1 = SM
210   B(I,JB) = (B(I,JB)-SM1)/A(I,I)
C
C      COMPLETE COMPUTATION OF SOLUTION VECTOR
C
      IF(K.EQ.N) GO TO 240
      DO 220 J = KP1,N
220   B(J,JB) = SZERO
      DO 230 I = 1,K
230   CALL H12(2,I,KP1,N,A(I,1),MDA,G(I),B(1,JB),1,MDB,1)
C
C      REORDER THE SOLUTION VECTOR TO COMPENSATE FOR THE COLUMN
C      INTERCHANGES.

```



```
C
24Ø DO 25Ø JJ = 1,LDIAG
      J = LDIAG + 1 - JJ
      IF(IP(J).EQ.J) GO TO 25Ø
      L = IP(J)
      TMP = B(L,JB)
      B(L,JB) = B(J,JB)
      B(J,JB) = TMP
25Ø CONTINUE
26Ø CONTINUE
C
C      THE SOLUTION VECTORS, X, ARE NOW IN THE FIRST N ROWS OF THE
C      ARRAY B(.)
C
27Ø KRANK = K
      RETURN
      END
```


ALGORITHM 560

JNF, An Algorithm for Numerical Computation of the Jordan Normal Form of a Complex Matrix [F2]

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Key Words and Phrases: Jordan normal form, canonical form, eigenvalues, eigenvectors, principal vectors, block diagonal form
CR Categories: 5.14
Language: Fortran

DESCRIPTION

1. Introduction

The routines given here are the actual Fortran implementation of the algorithm presented and discussed in [2]. We describe in detail how to use the Fortran subroutines and how to reach the results from a call. We also give some comments on the code that might be of value when implementing the subroutines on a particular machine. The subroutines have been checked with the PFORT verifier [4] and the notation from [2] is used.

2. The User-Written Routine DECIDE

DECIDE is a user-written subroutine which makes it possible for the user to change the grouping and/or the values of the numerical multiple eigenvalues.

This routine is useful when we have some information on the eigenvalues and their multiplicities in advance, and want the eigenvectors and the principal vectors for the given matrix. For instance, if we know of physical reasons why zero should be the only possible multiple eigenvalue and all others simple, then the contents of the parameters should be changed as indicated in Table I.

If the user does not want to influence the grouping or the values of the eigenvalues, it is enough to define a dummy subroutine, i.e.,

```
SUBROUTINE DECIDE (NM, N, NDEL, CSHTR, CSHTI, NBLOCK, HR, HI)
RETURN
END
```

3. The Rank Determination Process and the RDEFL Routine

The rank determination process is described in [2, Section 2.2 and Section 3, Step 6 of the algorithm]. According to these two alternatives we have two subroutines

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Table I. Example of Use of the Subroutine DECIDE

On entry N = 10 NBLOCK = 6		On exit NBLOCK = 7	
NDEL	CSHT	NDEL	CSHT
0	-1.5	0	-1.5
1	0.002	1	0.0
5	2.4	5	$\overline{H}(6, 6)$
7	3.5	6	$\overline{H}(7, 7)$
8	8.6	7	3.5
9	9.5	8	8.6
10		9	9.5
		10	

RDEFLL—one for the singular value decomposition strategy and one for the Kublanovskaya (RQ) decomposition strategy. Only the routine based on singular value decomposition is presented here.

In the RDEFLL routine we use the CSVD algorithm [1] with one change; namely, we sort the singular values in increasing instead of decreasing order as in the original program.

4. How to Reach Any of the Vectors

The eigenvectors and principal vectors are stored in the arrays ZR, ZI(NM, N). To find the vectors, we use the information stored in NBLOCK, NDEL, NDB, NDEFLL, and NXT (see the parameter list of subroutine JNF).

The column indices K in ZR and ZI for eigenvectors, principal vectors, and principal chains of an eigenvalue are listed below in Algol for-statement notation.

- Find the eigenvectors of the multiple eigenvalue
 $I = 1, \dots, \text{NBLOCK}$.
 $J := \text{NDB}[I]$;
for $K := \text{NDEL}[I] + 1$ **step** 1 **until** $\text{NDEFLL}[J + 1] - 1$ **do**;
- Find the principal vectors of grade P to the eigenvalue $I = 1, \dots, \text{NBLOCK}$.
 $J := \text{NDB}[I] + P - 1$;
for $K := \text{NDEFLL}[J]$ **step** 1 **until** $\text{NDEFLL}[J + 1] - 1$ **do**;
- Find the principal chain which ends with the eigenvector in column L.
 First:
 $(A - \text{EV}[L] * I) * Z[L] = 0$ (eigenvector)
 Then:
 $K1 := L$;
for $K := \text{NXT}[K1]$ **while** $K > 0$ **do**
begin
 Here:
 $(A - \text{EV}[K] * I) * Z[K] = \text{SUPD}[K1] * Z[K1]$; (principal vector)
 $K1 := K$
end;

(Here we use the fact that the diagonal elements may be different as in [2, eq. (4.2)]. To get chains of the nilpotent B, $\text{EV}[K]$ should be replaced by $\text{EV}[L]$.)

In Table II we list the values of NDEL, NDEFLL, NXT, NDB, and SUPD for the first test matrix in [2, Section 5].

5. Comments on the Fortran Code

Since we use the EISPACK routines CBAL, COMHES, COMLR2, and CBABK2 [5], we have their representation of complex arrays, i.e., the real parts are represented in one array with the name ending in R (e.g., HR, ZR) and the imaginary parts are represented in one array with the name ending in I (e.g., HI, ZI). This representation makes it possible to execute most of the computations in real arithmetic.

We use real arithmetic in all steps of the algorithm except Steps 4 and 7, i.e., the elimination processes, where we use some complex arithmetic.

Table II. Example of Output of Index Vectors

I	NDEL	NDB	NDEFL	NXT	SUPD
1	0	1	1	3	1.9903
2	4	3	3	4	48.080
3	9	6	5	0	0
4	10	0	7	0	0
5	0	0	9	7	3.1314
6	0	0	10	8	5.5354
7	0	0	11	0	0
8	0	0	0	9	1.3994
9	0	0	0	0	0
10	0	0	0	0	0

Note. Output for the first test matrix in [2] is displayed.

Because of this complex arithmetic we also have to use the complex standard functions `REAL(X)`, `AIMAG(X)`, and `CMPLX(X, Y)`.

If the actual Fortran compiler does not accept complex arithmetic, we can implement the functions `CMUL`, `CDIV`, `CADD`, and `CSUB` from [6] and substitute the actual statements with new statements.

A new parameter (`NOBACK`) has been added to the `EISPACK` routine `COMLR2`. We use this parameter (`NOBACK = 0`) as a control so that we do not get the eigenvectors from `COMLR2`. We get the transformations which transform the matrix to upper triangular form.

In the `RDEFL` routine we use complex arrays and arithmetic since the routine `CSVD1` [1] is made for complex arithmetic.

6. Numerical Experiments

We have performed several numerical tests, both on matrices with a well-defined Jordan normal form, and on cases specially constructed in order to provide the routine with difficulties. Results from full Jordan normal form reduction are reported and discussed in [2]. In [2, Section 4] we describe how to choose tolerance parameters in the grouping procedure and in the procedure for determining the structures inside the invariant subspaces corresponding to different numerical multiple eigenvalues. We also describe how to analyze the results from the program, given a combination of the tolerance parameters.

A large number of results from block diagonal reductions (`ISTEP = 5` and `6`) are reported and discussed in [3].

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ALGORITHM

[Summary information and a part of the listings are printed here. The complete listing is available from the ACM Algorithms Distribution Service (see inside back cover for order form) or may be found in microfiche form in "Collected Algorithms from ACM."]

`NAME(n)`: indicates a Fortran module with n records

NAME^T(n): indicates "NAME" is included for testing purposes
 Contents: MAIN^T(328), RESID^T(75), DECIDE^T(56), MATRIS^T(43),
 MATRS1^T(32), MATRS3^T(55), MATRS4^T(62), MATRS5^T(42),
 JNF(834), RDEFL(132), CBAL(209), COMHES(138),
 COMLR2(405), CBABK2(91), CVSD1(384), OUTPUT^T(33),
 MATBLK^T(37), MATRS2^T(148)

```

      SUBROUTINE DECIDE (NM,N,NDEL,CSHTR,CSHTI,NBLOCK,HR,HI)
C *****
C *
C * DECIDE IS A USER WRITTEN SUBROUTINE WHICH MAKES IT POSSIBLE
C * FOR THE USER TO CHANGE THE GROUPING OF THE NUMERICAL MUL-
C * TIPLE EIGENVALUES,AND/OR CHANGE THE VALUES OF THE MULTIPLE
C * EIGENVALUES.
C * THE FORMAL PARAMETER LIST
C *
C * NM      THE ROW DIMENSION OF THE TWO-DIMENSIONAL ARRAY
C *         PARAMETERS (HR,HI)AS DECLARED IN THE CALLING
C *         PROGRAM DIMENSION STATEMENT
C * N       THE ORDER OF THE ORIGINAL MATRIX
C * NBLOCK  THE NUMBER OF BLOCKS I.E THE NUMBER OF NUMERICAL
C *         MULTIPLE EIGENVALUES
C * CSHTR,  THE REAL AND IMAGINARY PARTS,RESPECTIVELY,OF THE
C * CSHTI   NUMERICAL MULTIPLE EIGENVALUES
C * NDEL    INDICATES THE MULTIPLICITY OF THE NUMERICAL
C *         MULTIPLE EIGENVALUES.NDEL(I+1)-NDEL(I)=THE
C *         MULTIPLICITY OF EIGENVALUE I FOR I=1,..,NBLOCK
C * HR,HI   THE REAL AND IMAGINARY PARTS,RESPECTIVELY, OF
C *         THE UPPER TRIANGULAR MATRIX RESULTING FROM
C *         STEPS 1-3 OF THE ALGORITHM
C *
C * NOTE -- THE USER MUST NOT CHANGE HR,HI
C *
C *****
      INTEGER NDEL(N)
      REAL HR(NM,N),HI(NM,N),CSHTR(1),CSHTI(1)
      WRITE(6,9)
9     FORMAT(1H0,1X,47HTHE FOLLOWING OUTPUT (A,B AND C) ARE PRINTED BY,
      *      31HTHE USER WRITTEN ROUTINE DECIDE/
      *      1H0,1X,31HSEE SECTION 2 OF THE ALGORITHM.)
      WRITE(6,10)
      WRITE(6,11) (HR(K,K),K=1,N)
      WRITE(6,11) (HI(K,K),K=1,N)
      WRITE(6,12)
      DO 1 K=1,NBLOCK
      MULT=NDEL(K+1)-NDEL(K)
      WRITE(6,13) NDEL(K+1),MULT,CSHTR(K),CSHTI(K)
1     CONTINUE
      WRITE(6,14)
14    FORMAT(1H0,1X,34HC--IN STEP 6 OF THE ALGORITHM THE ,
      *      50HSTRUCTURE OF EACH MULTIPLE EIGENVALUE IS COMPUTED.
      * /1H0,1X,42HFOR THAT REASON RDEFL SUCCESSIVELY COMPUTES,
      * 44H SINGULAR VALUE DECOMPOSITIONS. RDEFL PRINTS
      * /1H0,1X,39HTHE RESULTS BELOW(SEE ALSO COMMENTS IN ,
      * 35HRDEFL AND STEP 6 OF THE ALGORITHM.)
10   FORMAT(1H0,1X,43HA--ENTER DECIDE WITH EIGENVALUES COMPUTED ,
      *      36HBY COMLR2 (STEP 1 OF THE ALGORITHM ))
11   FORMAT(10F13.9)
12   FORMAT(1X,40HB--GROUPINGS OF THE EIGENVALUES,COMPUTED,
      *      27H BY STEP 3 OF THE ALGORITHM)
13  FORMAT(12H DIVISION AT,I4, 7H MULT.=,I4, 8H CENTER=,2E20.10)
      RETURN
      END
      SUBROUTINE JNF (NM,N,HR,HI,ZR,ZI,EVR,EVI,SUPD,NXT,NDEL,NDEFL,NDB,
C *****
C *
C * THE FORMAL PARAMETER LIST

```

```

C *                                00007020
C * ON INPUT -                    00007030
C *                                00007040
C * NM - MUST BE SET TO THE ROW DIMENSION OF THE TWO- 00007050
C * DIMENSIONAL ARRAY PARAMETERS AS DECLARED IN 00007060
C * THE CALLING PROGRAM 00007070
C * N - IS THE ORDER OF THE MATRICES 00007080
C *                                00007090
C * HR,HI - CONTAIN THE REAL AND IMAGINARY PARTS,RESPECTIVELY 00007100
C * OF THE COMPLEX MATRIX 00007110
C * EIN - IS A TOLERANCE PARAMETER,CORRESPONDING TO 00007120
C * PERTURBATIONS OF HR,HI, AND IS USED IN THE 00007130
C * GROUPING OF NUMERICAL MULTIPLE EIGENVALUES 00007140
C *                                00007150
C * TOL - IS A TOLERANCE PARAMETER USED IN THE CONSTRUCTION 00007160
C * OF THE NILPOTENT MATRICES (IS USED IN RDEFL) 00007170
C *                                00007180
C * ISTEP - RETURN FROM JNF AFTER STEP ISTEP(=1,2,3,4,5,6,7) 00007190
C * FULL REDUCTION TO JORDAN FORM IF ISTEP .GE. 7 00007200
C * OTHER CHOICES GIVE REDUCTION TO 00007210
C *                                00007220
C * 1 UPPER TRIANGULAR FORM 00007230
C *                                00007240
C * 2 UPPER TRIANGULAR FORM WITH THE EIGEN- 00007250
C * VALUES SORTED SUCH THAT CLOSE EIGEN- 00007260
C * VALUES APPEAR IN ADJACENT POSITIONS 00007270
C *                                00007280
C * 3 THE SAME AS 2,IN ADDITION GROUPING OF 00007290
C * CLOSE EIGENVALUES INTO BLOCKS CORRE- 00007300
C * SPONDING TO NUMERICAL MULTIPLE EIGEN- 00007310
C * VALUES 00007320
C *                                00007330
C * 4 BLOCK DIAGONAL UPPER TRIANGULAR FORM 00007340
C *                                00007350
C * 5 BLOCK DIAGONAL UPPER TRIANGULAR FORM 00007360
C * AND THE INVARIANT SUBSPACES CORRESPONDING 00007370
C * TO THE DIAGONALS BLOCKS HAVE UNITARY BASES 00007380
C *                                00007390
C * 6 THE STRUCTURE OF EACH DIAGONAL BLOCK IS 00007400
C * DETERMINED 00007410
C * ON OUTPUT - 00007420
C *                                00007430
C * THE PARAMETERS CONTAIN USEFUL INFORMATION DEPENDING ON 00007440
C * THE VALUE OF THE INPUT PARAMETER ISTEP. 00007450
C *                                00007460
C * NOTATION - (IF ISTEP .GE. N ) 00007470
C * WHERE N IS 1,2,3,4,5,6 OR 7 00007480
C *                                00007490
C *                                00007500
C * HR,HI - HAVE BEEN DESTROYED 00007510
C *                                00007520
C * IERR - IS A CONVERGENCE PARAMETER SET BY COMLR2 00007530
C * 0 FOR NORMAL RETURN 00007540
C * J IF THE J-TH EIGENVALUE HAS NOT BEEN DETERMINED 00007550
C * AFTER 30 ITERATIONS 00007560
C * IF AN ERROR EXIT IS MADE (IERR.NE.0),NONE 00007570
C * OF THE FOLLOWING PARAMETERS CONTAIN MEANINGFUL 00007580
C * RESULTS 00007590
C * (IF ISTEP.GE.1) 00007600
C * EVR,EVI CONTAIN THE REAL AND IMAGINARY PARTS,RESPECTIVELY 00007610
C * OF THE EIGENVALUES. 00007620
C * IF ISTEP.GE.1 BUT .LT.5 THE EIGENVALUES ARE 00007630
C * COMPUTED BY COMLR2. 00007640
C * IF ISTEP.GE.6, EVR,EVI CONTAIN THE MAIN DIAGONAL 00007650
C * OF THE JORDAN MATRIX. 00007660
C *                                00007670
C * NBLOCK- IS THE NUMBER OF NUMERICAL MULTIPLE EIGENVALUES 00007680
C * (IF ISTEP.GE.3) 00007690
C *                                00007700
C * NDEL -- INDICATES THE MULTIPLICITIES OF THE NUMERICAL 00007710
C * MULTIPLE EIGENVALUES 00007720
C * (IF ISTEP.GE.3) 00007730
C * NDEL(I+1)-NDEL(I)= MULTIPLICITY OF EIGENVALUE 00007740
C * I FOR I=1,...,NBLOCK 00007750
C * NDEFL - INDICATES THE STRUCTURE OF HR,HI 00007760
C * (IF ISTEP.GE.6) 00007770

```

```

C *
C * SUPD - CONTAINS THE REAL SUPERDIAGONAL OF THE JORDAN 00007780
C * MATRIX (THE COUPLING ELEMENTS) 00007790
C * (IF ISTEP.GE.7) 00007800
C * 00007810
C * 00007820
C * NXT - CONTAINS THE COLUMN INDICES OF SUPD 00007830
C * NXT(I)=J IMPLIES THAT THE VECTOR WITH COUPLING 00007840
C * ELEMENT SUPD(I) IS PLACED IN COLUMN 00007850
C * J OF ZR,ZI 00007860
C * (IF ISTEP.GE.7) 00007870
C * NDB - CONTAINS POINTER REFERENCES 00007880
C * NDB(I)=J MEANS THAT INFORMATION ABOUT THE 00007890
C * STRUCTURE OF THE NUMERICAL MULTIPLE 00007900
C * EIGENVALUE I STARTS IN POSITION J 00007910
C * OF NDEFL 00007920
C * (IF ISTEP.GE.6) 00007930
C * ZR,ZI - CONTAIN THE REAL AND IMAGINARY PARTS,RESPECTIVELY 00007940
C * OF THE ACCUMULATED TRANSFORMATIONS FROM STEP 1 00007950
C * TO 7 OF THE ALGORITHM. 00007960
C * IF ISTEP.GE.7, ZR,ZI ARE THE EIGENVECTORS AND 00007970
C * THE PRINCIPAL VECTORS. 00007980
C * 00007990
C * DELE - CONTAINS INFORMATION ABOUT DELETED ELEMENTS 00008000
C * DURING THE PROCESSES OF FINDING NILPOTENT MATRICES 00008010
C * DELE(I)= EUCLIDEAN NORM OF DELETED PART FOR BLOCK 00008020
C * I FOR I=1,...,NBLOCK 00008030
C * (IF ISTEP.GE.6) 00008040
C * 00008050
C * SM - CONTAINS ESTIMATES OF THE SPECTRAL PROJECTORS 00008060
C * CORRESPONDING TO DIFFERENT NUMERICAL MULTIPLE 00008070
C * EIGENVALUES 00008080
C * (IF ISTEP.GE.4) 00008090
C * 00008100
C * FOR COMPLETE DESCRIPTION SEE COMPUTATIONAL DETAILS IN 00008110
C * THE TEXT 00008120
C * 00008130

```


ALGORITHM 561

Fortran Implementation of Heap Programs for Efficient Table Maintenance [Z]

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of Commerce National Bureau of Standards

Key Words and Phrases: heap, table maintenance

CR Categories: 3.74, 5.31

Language: Fortran

DESCRIPTION

1. Introduction

In many computational problems one has a collection of entities whose time history is being studied. The calculation proceeds by creating, destroying, and generally processing these entities, henceforth called nodes. Often the node to be processed is that one which is "most important" in some problem-dependent sense. Examples of such calculations occur in particle following for accelerator problems and adaptive quadrature, to name but two [3, 4, 6]. In such problems the time spent maintaining the nodes is significant, and it is important to implement this efficiently.

An organization of nodes based on "most important in, first out" (i.e., where every deletion removes the most important node) is called a priority queue [5]. The computer implementation of such queues can vary considerably. An unordered table of nodes is the simplest arrangement but requires an expensive search for the most important node. Maintaining an ordered table allows easy access to the most important node, but then insertions become slow. A "heap" [1] is a data structure that implements a priority queue in an efficient way, being equally fast for both insertions and deletions. Here we present a suite of Fortran programs of library quality for heaps. These programs are of modest length and can easily be included in other packages. Similar programs were used in [6] but they were not found to be general enough for library use.

In this paper we do not describe any of the detailed properties of heaps since these have been well documented [1, 7]. Rather we give only enough information about them so that the use of the programs can be quickly understood. The last section contains two simple examples.

In our implementation the work required to form the heap from N given nodes is $O(N)$, and the work required either to remove the most important and update or to insert and update is $O(\log_2 N)$. This data structure is appropriate if the current most important node might be processed and destroyed before other

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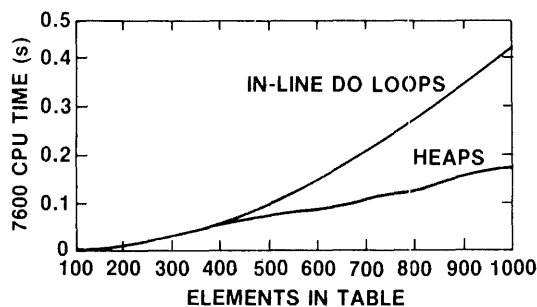


Fig. 1. Example of table management with heaps versus in-line DO loops.

nodes are created; the data structure is less appropriate if all the nodes are created before any are processed and destroyed.

The advantages of the programs presented in this paper over more familiar searching techniques are illustrated in Figure 1. Here we compare calls to the heap subroutines with in-line DO loops for the management of variously sized tables. This computation was selected to maximize the advantage of the in-line programming. (Specifically, example 1 of Section 3 was reprogrammed to use a DO loop at each stage to find the maximum element in an unordered table.) More typical problems should produce even more favorable time comparisons. For example, the production program MARLOWE, available from the Argonne Code Center, can have its execution time reduced 95 percent by using heaps instead of linear searches.

2. Using the Programs

We take a node to be one or more consecutive words of computer memory, with the data in these words defining the node. To investigate particle interactions (e.g., in an accelerator beam), each node could represent one particle, the nodal fields being position in phase space, energy, etc. For adaptive quadrature a node represents a portion of the region of integration; the fields in the node give the physical description of the region as well as local quadrature and error estimates. The most important node could be associated with the particle of largest energy or, for quadrature, the subregion with largest or smallest error, error per unit volume, that region nearest the origin, etc. To implement this variety, we require a Boolean function, HPFUN, on the nodes. Given two nodes A and B, this function returns .TRUE. if node A is "more important" than node B. The most important node is called TOP (i.e., it is the one for which HPFUN (TOP, X) = .TRUE. for any other node X).

Operations that can be performed on the heap, corresponding subroutines, and calling sequence parameters are shown in Table I. We also show (for completeness) the form of the Boolean HPFUN as well as an internally used subroutine HPGRO. A description of the parameters is given later.

The usual sequence of calls begins with HPINIT, perhaps followed by HPBLD. The main calculations usually require numerous calls to HPINS, HPACC, and HPDEL in various orders and combinations.

The package stores nodes into and removes them from a scratch area (DATA) provided by the user. The function HPFUN is an argument to the subroutines and thus may be changed from run to run or even during a single run. The latter situation might be useful during a calculation in which too many nodes are being generated.

The subroutine parameters are as follows:

NMAX	Maximum number of nodes permitted, set by user.
NWDS	Number of words per node, set by user.
DATA	Real work array, at least NMAX * NWDS words, dimensioned by the user for node storage.

Table I. Heap Functions, Subroutines, and Calling Sequence Parameters

Function	Subroutine	Calling sequence parameters
Initialize package at start of each calculation	HPINIT	NMAX, N, T
Build a heap	HPBLD	NMAX, NWDS, DATA, N, T, HPFUN
Insert a node	HPINS	NMAX, NWDS, DATA, N, T, XNODE, HPFUN
Delete top node	HPDEL	NMAX, NWDS, DATA, N, T, HPFUN
Access a node ($K = 1$ gives top)	HPACC	NMAX, NWDS, DATA, N, T, XNODE, K
Determine the "more important" of two given nodes	HPFUN ^a	A, B, NWDS
Internally used subroutine	HPGRO	NMAX, NWDS, DATA, N, T, HPFUN, I

^a Generic name—user may supply his own, but name must be EXTERNALed.

- N** Current number of nodes in the heap, set by the package.
- T** Integer work array dimensioned at least NMAX by the user, defined and used internally as a pointer.
- XNODE** Real array of NWDS words dimensioned by the user and used to store a single node. Thus to insert a node via HPINS, the user forms XNODE, i.e., defines its NWDS words and calls HPINS to copy these into the heap. To access a node with HPACC, the package copies the nodal data into XNODE (see examples).
- HPFUN** Name of a logical function with two real parameters and one integer parameter, of the form

```
LOGICAL FUNCTION HPFUN (A, B, NWDS)
INTEGER NWDS
REAL A(NWDS), B(NWDS)
```

The name of this function must be included in an EXTERNAL statement in the user calling program. When this function is called by the package, the arguments A and B will be two different nodes. The package provides two Boolean functions, LESS and GREATR, defined by

```
LESS = A(1) .LT. B(1)
GREATR = A(1) .GT. B(1)
```

These functions may be used when the characteristic on which the heap is organized appears in the first word of each node. The LESS and GREATR functions make the TOP smallest or largest, respectively. If the user wants more complicated functions, he must write them himself. For example, to select the node that is closest to the origin in the sense of the sum of squares of components, we can write

```
LOGICAL FUNCTION L2 (A, B, NWDS)
INTEGER NWDS, I
REAL A (NWDS), B (NWDS), DA, DB
DA = 0.0
DB = 0.0
DO 1 I = 1, NWDS
1  DA = DA + A (I)**2
   DB = DB + B (I)**2
L2 = DA .LT. DB
RETURN
END
```

K (In SUBROUTINE HPACC) the index of the node to be accessed; set by user. $K = 1$ corresponds to the TOP, and K must be in the range $1 \leq K \leq N \leq NMAX$. (Since the heap is only partially ordered, there is no way to predict the node produced by using this subroutine with $K > 1$. It is allowed as a convenience to those users who may wish to examine all the nodes without altering the heap structure.)

Note. If subroutines in the package are called with incorrect arguments, e.g., $N > NMAX$, the programs return without performing any calculations. No error diagnostics are given.

Operation Counts. The initialization program HPINIT and the infrequently used HPBLD require $O(N)$ operations. Inserting or deleting a node, HPINS or HPDEL, requires $O(\log_2 N)$ operations; accessing a node with HPACC requires $O(1)$ operations and has no effect upon the heap.

3. Examples of Use

Our first example is trivial. It has no application except to illustrate the use of the package. Nevertheless it forms the framework for the second example, adaptive quadrature, an important algorithm.

Example 1. A Simple Example of Heap Subroutines

A node is a single real number, between 0 and 1, generated at random. The TOP node in the heap is that one with largest numerical value. The calculation is as follows. Start with an initial node. While the number of nodes is less than 1000, print the value of the TOP node, remove it from the heap, and replace it by two others also selected at random. The complete Fortran program is as follows:

```

      REAL DATA (1000), XNODE
      INTEGER T(1000)
      EXTERNAL GREATR
      LOGICAL GREATR
C
      NMAX = 1000
      CALL HPINIT (NMAX, N, T)
      XNODE = RANNUM (0.)
      CALL HPINS (NMAX, 1, DATA, N, T, XNODE, GREATR)
C
1  IF (N .GE. NMAX) STOP
      CALL HPACC (NMAX, 1, DATA, N, T, XNODE, 1)
      WRITE (6,100) XNODE
      CALL HPDEL (NMAX, 1, DATA, N, T, GREATR)
      XNODE = RANNUM (0.)
      CALL HPINS (NMAX, 1, DATA, N, T, XNODE, GREATR)
      XNODE = RANNUM (0.)
      CALL HPINS (NMAX, 1, DATA, N, T, XNODE, GREATR)
      GO TO 1
100  FORMAT (E16.8)
      END

```

Example 2. Global Adaptive Quadrature

We start with an initial interval, a user-provided subroutine for evaluating the integrand, and $\epsilon > 0$. The algorithm generates subintervals by bisection and produces an integral estimate and error estimate for each. The subinterval subdivided next is that one having largest error estimate. The algorithm terminates where the sum of subinterval error estimates is less than ϵ . This is the prototype algorithm on which heaps were successfully used in [6]. A node is a real four-word array corresponding to an interval and has the form

$$\begin{aligned} \text{NODE}_1 &= \text{error estimate} \geq 0 \\ \text{NODE}_2 &= \text{integral estimate} \end{aligned}$$

NODE₃ = left endpoint
 NODE₄ = right endpoint

The heap is organized so that the TOP is the node with the largest error estimate. Thus the function GREATR provided in the package is used.

The main user subroutine is AGQ. Auxiliary calculations use the user-written programs:

QINIT Initializes the first interval to a node.
 QSDVD Produces two half-interval nodes from each full interval.
 QUAD Provides integral estimate and error estimate for the user function on a given node.

In our sample program integral/error estimates are calculated by 2- and 3-point Gauss quadrature. This is, however, independent of any other aspect of the algorithm, and readers may substitute Newton-Cotes, Kronrod, etc. Our program has not taken advantage of the fact that for this problem a deletion is always followed by an insertion; the gain by doing this is very small. Furthermore the favorable timings of the heap programs relative to a simple list only appear after about 100 nodes, a situation which does not occur for the two problems in the example. On the other hand, for automatic integrators small gains in efficiency are not a major issue for problems generating such a small number of nodes. This same simple program organization can be used, however, for multidimensional adaptive quadrature after obvious changes in the definition and storage requirements for a node. For these problems large numbers of nodes can easily be generated. The program of Example 2 is a prototype of the library software that is now being developed in this area. For more complex applications of these ideas see [2] and [4].

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ALGORITHM

```

REAL WORK(1000), TNODE(4)                                MAIN 1
INTEGER IWORK(200)                                       MAIN 2
EXTERNAL F,G,LEFT                                        MAIN 3
LOGICAL LEFT                                             MAIN 4
C                                                         MAIN 5
C     TEST DRIVER TO ILLUSTRATE THE USE OF HEAP SUBROUTINES FOR MAIN 6
C     QUADRATURE.                                         MAIN 7
C     THIS PROGRAM CALLS THE ADAPTIVE QUADRATURE SUBROUTINE AGQ MAIN 8
C     TO INTEGRATE THE TWO FUNCTIONS                     MAIN 9
C     F(X)=EXP(X)                                         MAIN 10
C     AND                                                  MAIN 11
C     G(X)=1./SQRT(X)                                     MAIN 12
C     ON (0., 1.) TO AN ABSOLUTE ACCURACY OF 1.E-05     MAIN 13
C     THE DEFINITIONS OF THE ARGUMENTS TO SUBROUTINE AGQ ARE GIVEN MAIN 14
C     IN THE COMMENTS FOR THAT SUBROUTINE.             MAIN 15
C     FOR EACH OF THESE TWO PROBLEMS THIS TEST PROGRAM PRINTS MAIN 16

```

```

C          Q = THE COMPUTED ESTIMATE OF THE INTEGRAL          MAIN 17
C          E = THE COMPUTED ESTIMATE OF THE ERROR            MAIN 18
C          N = THE NUMBER OF SUBINTERVALS IN THE FINAL PARTITION OF MAIN 19
C              THE INTERVAL (0., 1.)                          MAIN 20
C          IFLAG = A SUCCESS (0) OR FAILURE (1) INDICATOR.    MAIN 21
C                                                            MAIN 22
C          CALL AGQ(0., 1., F, 1.E-05, Q, E, N, 100, WORK, IWORK, IFLAG) MAIN 23
C          WRITE (6, 102)                                     MAIN 24
C          WRITE(6, 100) Q, E, N, IFLAG                       MAIN 25
C          CALL AGQ(0., 1., G, 1.E-05, Q, E, N, 100, WORK, IWORK, IFLAG) MAIN 26
C          WRITE(6, 100) Q, E, N, IFLAG                       MAIN 27
C                                                            MAIN 28
C          AFTER THE SECOND PROBLEM WE REORGANIZE THE HEAP SO THAT THE MAIN 29
C              ROOT INTERVAL IS THE LEFTMOST ONE, AND THEN PRINT THE MAIN 30
C              SUBINTERVALS USED LEFT TO RIGHT.              MAIN 31
C                                                            MAIN 32
C          CALL HPBLD(100, 4, WORK, N, IWORK, LEFT)           MAIN 33
C          M=N                                                 MAIN 34
C          WRITE(6, 103)                                       MAIN 35
C          DO 1 I=1, M                                         MAIN 36
C              CALL HPACC(100, 4, WORK, N, IWORK, TNODE, 1)   MAIN 37
C              WRITE(6, 101) (TNODE(J), J=1, 4)               MAIN 38
C              CALL HPDEL(100, 4, WORK, N, IWORK, LEFT)       MAIN 39
C          1 CONTINUE                                         MAIN 40
C                                                            MAIN 41
C          100 FORMAT(1X, 2E16.8, 2I10)                         MAIN 42
C          101 FORMAT( 1X, 4E16.8)                             MAIN 43
C          102 FORMAT(16H  QUAD. EST.  ,16H  ERR. EST.  ,10H NO. INTS ,10H SMAIN 44
C              1SUCCESS=0 )                                    MAIN 45
C          103 FORMAT(65H FINAL SUBINTERVALS...ERR EST, QUAD EST, L. END PT, RTMAIN 46
C              1. END PT  )                                    MAIN 47
C          STOP                                                 MAIN 48
C          END                                                 MAIN 49

          LOGICAL FUNCTION LEFT(A, B, NCHAR)                    LEFT 1
          DIMENSION A(1), B(1)                                  LEFT 2
C
C          A TEST HEAP FUNCTION.                                LEFT 3
C          THIS FUNCTION SETS LEFT = TRUE IF THE THIRD COMPONENT OF LEFT 4
C              NODE A IS .LT. THE THIRD COMPONENT OF NODE B.  LEFT 5
C          FOR THE TEST PROBLEM THIS COMPONENT CONTAINS THE LEFT END- LEFT 6
C              POINT OF THE INTERVAL DEFINED BY THE NODE.     LEFT 7
C                                                            LEFT 8
C          LEFT= A(3) .LT. B(3)                                 LEFT 9
C          RETURN                                               LEFT 10
C          END                                                  LEFT 12

          FUNCTION F(X)                                         F 1
          F=EXP(X)                                              F 2
          RETURN                                               F 3
          END                                                  F 4

          FUNCTION G(X)                                         G 1
          G=1./SQRT(X)                                          G 2
          RETURN                                               G 3
          END                                                  G 4

          SUBROUTINE AGQ(A, B, F, EPS, Q, E, N, NMAX, XNODES, T, IFLAG) AGQ 1
C
C          1-D ADAPTIVE QUADRATURE-EXAMPLE USING HEAPS          AGQ 2
C          INPUT                                                AGQ 3
C          A, B = LIMITS OF INTEGRATION A.LT.B                 AGQ 4
C          F = NAME OF USER SUPPLIED INTEGRAND FUNCTION, FUNCTION F(X) AGQ 5
C          EPS = REQUESTED ABSOLUTE ACCURACY OF QUADRATURE     AGQ 6
C          NMAX = MAX NUMBER OF NODES ALLOWED                  AGQ 7
C          XNODES = ARRAY FOR NODE STORAGE, DIMENSIONED AT LEAST 4*NMAX AGQ 8
C          T = INTEGER ARRAY USED INTERNALLY, DIMENSIONED AT LEAST NMAX AGQ 9
C          OUTPUT                                               AGQ 10
C          Q = QUADRATURE ESTIMATE                              AGQ 11
C          E = ERROR ESTIMATE                                    AGQ 12
C          N = NUMBER OF NODES REQUIRED FOR THE QUADRATURE       AGQ 13
C

```

```

C      IFLAG = ERROR INDICATOR                      AGQ 15
C      = 0 PROGRAM APPARENTLY SUCCESSFUL           AGQ 16
C      = 1 NODE STORAGE EXCEEDED, CALCULATION TERMINATED AGQ 17
C
C      EACH NODE IN THE HEAP REPRESENTS A SUBINTERVAL AND IS AGQ 18
C      A FOUR WORD ARRAY.                          AGQ 19
C      NODE(1) = ERROR ESTIMATE                    AGQ 20
C      NODE(2) = QUADRATURE ESTIMATE               AGQ 21
C      NODE(3) = LEFT INTERVAL ENDPOINT            AGQ 22
C      NODE(4) = RIGHT INTERVAL ENDPOINT           AGQ 23
C
C      EXTERNAL GREATR, F                          AGQ 24
C      LOGICAL GREATR                             AGQ 25
C      REAL TNODE(4), N1(4), N2(4), XNODES(1)     AGQ 26
C      INTEGER T(1)                                AGQ 27
C
C      CALL HPINIT(NMAX, N, T)                     AGQ 28
C      CALL QINIT(TNODE, A, B)                    AGQ 29
C      E=TNODE(1)                                  AGQ 30
C      Q=TNODE(2)                                  AGQ 31
C
C      INSERT INITIAL NODE                         AGQ 32
C
C      CALL HPINS(NMAX, 4, XNODES, N, T, TNODE, GREATR) AGQ 33
C
C      I IF(E .LE. EPS) RETURN                    AGQ 34
C      CALL HPACC(NMAX, 4, XNODES, N, T, TNODE, 1) AGQ 35
C      CALL HPDEL(NMAX, 4, XNODES, N, T, GREATR)  AGQ 36
C      CALL QSDVD(TNODE, N1, N2, F)               AGQ 37
C      CALL HPINS(NMAX, 4, XNODES, N, T, N1, GREATR) AGQ 38
C      CALL HPINS(NMAX, 4, XNODES, N, T, N2, GREATR) AGQ 39
C      E=E-TNODE(1)+N1(1)+N2(1)                  AGQ 40
C      Q=Q-TNODE(2)+N1(2)+N2(2)                  AGQ 41
C      IFLAG=1                                     AGQ 42
C      IF(N .GE. NMAX-1) RETURN                  AGQ 43
C      IFLAG=0                                     AGQ 44
C      GO TO 1                                     AGQ 45
C      END                                         AGQ 46
C
C      SUBROUTINE QINIT(NODE, A, B)                QINIT 1
C      INITIALIZES THE FIRST INTERVAL TO A NODE BY SETTING THE QINIT 2
C      ERROR ESTIMATE AND QUADRATURE ESTIMATE TO 1000. AND 0. RESP. QINIT 3
C
C      REAL NODE(4)                                QINIT 4
C      NODE(1)=1000.                               QINIT 5
C      NODE(2)=0.0                                QINIT 6
C      NODE(3)=A                                   QINIT 7
C      NODE(4)=B                                   QINIT 8
C      RETURN                                       QINIT 9
C      END                                         QINIT 10
C
C      SUBROUTINE QSDVD(NODE, R, L, F)             QSDVD 1
C      SUBDIVIDES A NODE INTO TWO EQUAL HALVES.   QSDVD 2
C      THE ENDPOINTS OF THE NEW NODES ARE COMPUTED HERE. THE ERROR QSDVD 3
C      AND QUADRATURE ESTIMATES FOR THE LEFT (L) AND RIGHT (R) QSDVD 4
C      HALVES ARE COMPUTED BY THE SUBROUTINE QUAD. QSDVD 5
C
C      EXTERNAL F                                  QSDVD 6
C      REAL NODE(4), R(4), L(4)                   QSDVD 7
C      L(3)=NODE(3)                               QSDVD 8
C      R(4)=NODE(4)                               QSDVD 9
C      L(4)=(L(3)+R(4))/2.                        QSDVD 10
C      R(3)=L(4)                                  QSDVD 11
C      CALL QUAD(L, F)                             QSDVD 12
C      CALL QUAD(R, F)                             QSDVD 13
C      RETURN                                       QSDVD 14
C      END                                         QSDVD 15
C
C      SUBROUTINE QUAD(NODE, F)                   QUAD 1
C      ESTIMATES INTEGRAL OF F ON NODE USING 2 AND 3 POINT GAUSS. QUAD 2
C

```

```

C          THE 3 POINT FORMULA GIVES ESTIMATE AND DIFFERENCE OF 2 AND 3 QUAD 4
C          GIVES ERROR ESTIMATE QUAD 5
C QUAD 6
REAL NODE(4) QUAD 7
XL=(NODE(4)-NODE(3))/2. QUAD 8
XM=(NODE(4)+NODE(3))/2. QUAD 9
Q2=XL*(F(XL*(-.5773502692)+XM)+F(XL*(.5773502692)+XM)) QUAD 10
Q3=XL*((F(XL*(-.7745966692)+XM)+F(XL*(.7745966692)+XM))/1.8+ QUAD 11
1 F(XM)/1.125) QUAD 12
NODE(1)=ABS(Q2-Q3) QUAD 13
NODE(2)=Q3 QUAD 14
RETURN QUAD 15
END QUAD 16

C          H E A P PACKAGE HPINIT 1
C          A COLLECTION OF PROGRAMS WHICH MAINTAIN A HEAP DATA HPINIT 2
C          STRUCTURE. BY CALLING THESE SUBROUTINES IT IS POSSIBLE TO HPINIT 3
C          INSERT, DELETE, AND ACCESS AN EXISTING HEAP OR TO BUILD A HPINIT 4
C          NEW HEAP FROM AN UNORDERED COLLECTION OF NODES. THE HEAP HPINIT 5
C          FUNCTION IS AN ARGUMENT TO THE SUBROUTINES ALLOWING VERY HPINIT 6
C          GENERAL ORGANIZATIONS. HPINIT 7
C          THE USER MUST DECIDE ON THE MAXIMUM NUMBER OF NODES HPINIT 8
C          ALLOWED AND DIMENSION THE REAL ARRAY DATA AND THE INTEGER HPINIT 9
C          ARRAY T USED INTERNALLY BY THE PACKAGE. THESE VARIABLES ARE HPINIT10
C          THEN PASSED THROUGH THE CALL SEQUENCE BETWEEN THE HEAP HPINIT11
C          PROGRAMS BUT ARE NOT IN GENERAL ACCESSED BY THE USER. HE HPINIT12
C          MUST ALSO PROVIDE A HEAP FUNCTION WHOSE NAME MUST BE INCLUD- HPINIT13
C          ED IN AN EXTERNAL STATEMENT IN THE USER PROGRAM WHICH CALLS HPINIT14
C          THE HEAP SUBROUTINES. TWO SIMPLE HEAP FUNCTIONS ARE HPINIT15
C          PROVIDED WITH THE PACKAGE. HPINIT16
C HPINIT17
SUBROUTINE HPINIT(NMAX,N,T) HPINIT18
C HPINIT19
C PURPOSE HPINIT20
C THIS ROUTINE INITIALIZES THE HEAP PROGRAMS. HPINIT21
C IT IS CALLED ONCE AT THE START OF EACH NEW CALCULATION HPINIT22
C INPUT HPINIT23
C NMAX = MAXIMUM NUMBER OF NODES ALLOWED BY USER. HPINIT24
C OUTPUT HPINIT25
C N = CURRENT NUMBER OF NODES IN HEAP = 0. HPINIT26
C T = INTEGER ARRAY OF POINTERS TO POTENTIAL HEAP NODES. HPINIT27
C HPINIT28
INTEGER T(1) HPINIT29
DO I I = 1, NMAX HPINIT30
1 T(I)=I HPINIT31
N = 0 HPINIT32
RETURN HPINIT33
END HPINIT34

C          SUBROUTINE HPINS(NMAX,NCHAR,DATA,N,T,NODE,HPFUN) HPINS 1
C          HPINS 2
C PURPOSE HPINS 3
C THIS ROUTINE INSERTS A NODE INTO AN ALREADY EXISTING HEAP. HPINS 4
C THE RESULTING TREE IS RE-HEAPED. HPINS 5
C HPINS 6
C INPUT HPINS 7
C NMAX = MAXIMUM NUMBER OF NODES ALLOWED BY USER. HPINS 8
C NCHAR = NUMBER OF WORDS PER NODE. HPINS 9
C DATA = WORK AREA FOR STORING NODES. HPINS 10
C N = CURRENT NUMBER OF NODES IN THE TREE. HPINS 11
C T = INTEGER ARRAY OF POINTERS TO HEAP NODES. HPINS 12
C NODE = A REAL ARRAY, NCHAR WORDS LONG, WHICH HPINS 13
C CONTAINS THE NODAL INFORMATION TO BE INSERTED. HPINS 14
C HPFUN = NAME OF USER WRITTEN FUNCTION TO DETERMINE HPINS 15
C THE ROOT NODE. HPINS 16
C OUTPUT HPINS 17
C DATA = WORK AREA WITH NEW NODE INSERTED. HPINS 18
C N = UPDATED NUMBER OF NODES. HPINS 19
C T = UPDATED INTEGER POINTER ARRAY. HPINS 20
C HPINS 21
DIMENSION T(1),NODE(1),DATA(1) HPINS 22
REAL NODE HPINS 23
INTEGER T HPINS 24
LOGICAL HPFUN HPINS 25

```


IF(N .EQ. NMAX) RETURN	HPINS 26
N=N+1	HPINS 27
J= (T(N)-1)*NCHAR	HPINS 28
DO 1 I= 1,NCHAR	HPINS 29
IPJ=I+J	HPINS 30
1 DATA(IPJ) = NODE(I)	HPINS 31
J=N	HPINS 32
2 CONTINUE	HPINS 33
IF(J .EQ. 1) RETURN	HPINS 34
JT=T(J)	HPINS 35
J2=J/2	HPINS 36
JT2=T(J2)	HPINS 37
JL=(JT2-1)*NCHAR+1	HPINS 38
JR=(JT-1)*NCHAR+1	HPINS 39
IF(HPFUN(DATA(JL), DATA(JR), NCHAR)) RETURN	HPINS 40
T(J2)=T(J)	HPINS 41
T(J)=JT2	HPINS 42
J=J/2	HPINS 43
GO TO 2	HPINS 44
END	HPINS 45
SUBROUTINE HPBLD(NMAX, NCHAR, DATA, NELTS, T, HPFUN)	HPBLD 1
C	HPBLD 2
C PURPOSE	HPBLD 3
C BUILDS A HEAP, IN T, FROM AN ARRAY OF NELTS ELEMENTS	HPBLD 4
C IN DATA, WHICH ARE SPACED NCHAR APART.	HPBLD 5
C AT CONCLUSION OF CALCULATION THE ROOT SATISFIES	HPBLD 6
C HPFUN(ROOT, SON) = .TRUE. FOR ANY SON.	HPBLD 7
C USES SUBROUTINE HPGRO BY FEEDING IT ONE ELEMENT OF	HPBLD 8
C THE ARRAY AT A TIME.	HPBLD 9
C	HPBLD 10
C INPUT	HPBLD 11
C NMAX = MAXIMUM NUMBER OF NODES ALLOWED BY USER.	HPBLD 12
C NCHAR = NUMBER OF WORDS PER NODE.	HPBLD 13
C DATA = WORK AREA IN WHICH THE NODES ARE STORED.	HPBLD 14
C NELTS = CURRENT NUMBER OF NODES.	HPBLD 15
C T = INTEGER ARRAY OF POINTERS TO HEAP NODES.	HPBLD 16
C HPFUN = NAME OF USER WRITTEN FUNCTION TO DETERMINE ROOT NODE.	HPBLD 17
C OUTPUT	HPBLD 18
C DATA = WORK AREA IN WHICH THE NODES ARE STORED.	HPBLD 19
C T = INTEGER ARRAY OF POINTERS TO HEAP NODES.	HPBLD 20
C IN PARTICULAR T(1) POINTS TO THE ROOT.	HPBLD 21
C	HPBLD 22
C EXTERNAL HPFUN	HPBLD 23
C LOGICAL HPFUN	HPBLD 24
C DIMENSION DATA(1)	HPBLD 25
C INTEGER T(1)	HPBLD 26
C IF(NMAX .LT. NELTS) RETURN	HPBLD 27
C INDEX = NELTS/2	HPBLD 28
C 1 CONTINUE	HPBLD 29
C IF(INDEX .EQ. 0) RETURN	HPBLD 30
C CALL HPGRO(NMAX, NCHAR, DATA, NELTS, T, HPFUN, INDEX)	HPBLD 31
C INDEX = INDEX-1	HPBLD 32
C GO TO 1	HPBLD 33
C END	HPBLD 34
SUBROUTINE HPDEL(NMAX, NCHAR, DATA, NCELLS, T, HPFUN)	HPDEL 1
C	HPDEL 2
C PURPOSE	HPDEL 3
C DELETE ROOT ELEMENT OF HEAP. RESULTING TREE IS REHEAPED.	HPDEL 4
C INPUT	HPDEL 5
C NMAX = MAXIMUM NUMBER OF NODES ALLOWED BY USER.	HPDEL 6
C NCHAR = NUMBER OF WORDS PER NODE.	HPDEL 7
C DATA = WORK AREA IN WHICH THE NODES ARE STORED.	HPDEL 8
C NCELLS = CURRENT NUMBER OF NODES.	HPDEL 9
C T = INTEGER ARRAY OF POINTERS TO NODES.	HPDEL 10
C HPFUN = NAME OF USER WRITTEN FUNCTION TO DETERMINE ROOT NODE.	HPDEL 11
C OUTPUT	HPDEL 12
C NCELLS = UPDATED NUMBER OF NODES.	HPDEL 13
C T = UPDATED INTEGER POINTER ARRAY TO NODES.	HPDEL 14
C	HPDEL 15
C EXTERNAL HPFUN	HPDEL 16
C LOGICAL HPFUN	HPDEL 17

	DIMENSION DATA(1)	HPDEL 18
	INTEGER T(1)	HPDEL 19
	IF(NCELLS .EQ. 0) RETURN	HPDEL 20
	JUNK=T(1)	HPDEL 21
	T(1)=T(NCELLS)	HPDEL 22
	T(NCELLS)=JUNK	HPDEL 23
	NCELLS=NCELLS-1	HPDEL 24
	CALL HPGRO(NMAX,NCHAR,DATA,NCELLS,T,HPFUN,1)	HPDEL 25
	RETURN	HPDEL 26
	END	HPDEL 27
	SUBROUTINE HPACC(NMAX,NCHAR,DATA,N,T,NODE,K)	HPACC 1
C		HPACC 2
C	PURPOSE	HPACC 3
C	TO ACCESS THE K-TH NODE OF THE HEAP,	HPACC 4
C	1 .LE. K .LE. N .LE. NMAX	HPACC 5
C	INPUT	HPACC 6
C	NMAX = MAXIMUM NUMBER OF NODES ALLOWED BY USER.	HPACC 7
C	DATA = WORK AREA FOR STORING NODES.	HPACC 8
C	N = CURRENT NUMBER OF NODES IN THE HEAP.	HPACC 9
C	T = INTEGER ARRAY OF POINTERS TO HEAP NODES.	HPACC 10
C	NODE = A REAL ARRAY, NCHAR WORDS LONG, IN WHICH NODAL IN-	HPACC 11
C	FORMATION WILL BE INSERTED.	HPACC 12
C	K = THE INDEX OF THE NODE TO BE FOUND AND INSERTED INTO NODE.	HPACC 13
C		HPACC 14
C	OUTPUT	HPACC 15
C	NODE = A REAL ARRAY. CONTAINS IN NODE(1),...,NODE(NCHAR)	HPACC 16
C	THE ELEMENTS OF THE K-TH NODE.	HPACC 17
C		HPACC 18
	REAL DATA(1), NODE(1)	HPACC 19
	INTEGER T(1)	HPACC 20
	IF (K .LT. 1 .OR. K .GT. N .OR. N .GT. NMAX) RETURN	HPACC 21
	J=(T(K)-1)*NCHAR	HPACC 22
	DO 1 I=1,NCHAR	HPACC 23
	IPJ=I+J	HPACC 24
1	NODE(I)=DATA(IPJ)	HPACC 25
	RETURN	HPACC 26
	END	HPACC 27
	LOGICAL FUNCTION GREATR(A,B,NCHAR)	GREATR 1
	REAL A(1),B(1)	GREATR 2
	GREATR= A(1) .GT. B(1)	GREATR 3
	RETURN	GREATR 4
	END	GREATR 5
	LOGICAL FUNCTION LESS(A,B,NCHAR)	LESS 1
	REAL A(1),B(1)	LESS 2
	LESS= A(1) .LT. B(1)	LESS 3
	RETURN	LESS 4
	END	LESS 5
	SUBROUTINE HPGRO(NMAX,KD,DATA,NELTS,T,HPFUN,KTEMP)	HPGRO 1
C		HPGRO 2
C	PURPOSE	HPGRO 3
C	FORMS A HEAP OUT OF A TREE. USED PRIVATELY BY HPBLD.	HPGRO 4
C	THE ROOT OF THE TREE IS STORED IN LOCATION T(KTEMP).	HPGRO 5
C	FIRST SON IS IN LOCATION T(2KTEMP), NEXT SON	HPGRO 6
C	IS IN LOCATION T(2KTEMP+1).	HPGRO 7
C	THIS PROGRAM ASSUMES EACH BRANCH OF THE TREE IS A HEAP.	HPGRO 8
C		HPGRO 9
	DIMENSION T(1),DATA(1)	HPGRO 10
	INTEGER T	HPGRO 11
	LOGICAL HPFUN	HPGRO 12
	IF(NELTS .GT. NMAX) RETURN	HPGRO 13
C		HPGRO 14
	K=KTEMP	HPGRO 15
1	I=2*K	HPGRO 16
C		HPGRO 17
C	TEST IF ELEMENT IN I TH POSITION IS A LEAF.	HPGRO 18
C		HPGRO 19
	IF(I .GT. NELTS) RETURN	HPGRO 20

C		HPFRO	21
C	IF THERE IS MORE THAN ONE SON, FIND WHICH SON IS SMALLEST.	HPGRO	22
C		HPGRO	23
	IF(I .EQ. NELTS) GO TO 2	HPGRO	24
	ITEMP=T(I)	HPGRO	25
	ITP1=T(I+1)	HPGRO	26
	IL=(ITP1-1)*KD+1	HPGRO	27
	IR=(ITEMP-1)*KD+1	HPGRO	28
	IF(HPFUN(DATA(IL),DATA(IR),KD)) I=I+1	HPGRO	29
C		HPGRO	30
C	IF A SON IS LARGER THAN FATHER, INTERCHANGE	HPGRO	31
C	THIS DESTROYS HEAP PROPERTY, SO MUST RE-HEAP REMAINING	HPGRO	32
C	ELEMENTS	HPGRO	33
C		HPGRO	34
2	CONTINUE	HPGRO	35
	KT=T(K)	HPGRO	36
	ITEMP=T(I)	HPGRO	37
	IL=(KT-1)*KD+1	HPGRO	38
	IR=(ITEMP-1)*KD+1	HPGRO	39
	IF(HPFUN(DATA(IL),DATA(IR),KD)) RETURN	HPGRO	40
	ITEMP=T(I)	HPGRO	41
	T(I)=T(K)	HPGRO	42
	T(K)=ITEMP	HPGRO	43
	K=I	HPGRO	44
	GO TO 1	HPGRO	45
	END	HPGRO	46

ALGORITHM 562

Shortest Path Lengths [H]

U. PAPE
Technische Universität Berlin

Key Words and Phrases: shortest path, shortest route problem
CR Categories: 5.32
Language: Fortran IV

DESCRIPTION

This algorithm finds the shortest path lengths from a specific node to all other nodes in a network. It is a modification of Moore's algorithm [4], originally due to d'Esopo as reported by Pollack and Wiebenson [7] and refined by Pape [5, 6]. The algorithm does not determine the shortest paths. If a shortest path is to be determined, this algorithm may be combined with any algorithm that traces the shortest path from the root node to any other node from the list of predecessors of nodes. As the calculation of the path lengths is the most time-consuming part of the whole process, we confine our attention to this problem.

The main idea of the given algorithm centers around the "status" of a new-found successor node j . If j has not yet been reached by the process, it is entered at the end of the successor list; if it is currently in the list, no new entry is made, but if it has already been processed and removed from the list, it is entered at the top of the list so that it will be processed next. The current status of each node is recorded by the values of the list which are stored as linked deque (double-ended queue). The steps involved are as follows:

```

top element of the deque :=  $j$ ;
 $m_j := \infty$ ;  $m_{j_j} := 0$ ;
while deque  $\neq \emptyset$  do
  begin  $i :=$  top element of the deque;
    remove top element;
    for all successors  $k$  of  $i$  do
      begin  $m_{jk} := m_j + d_{ik}$ 
        if  $m_{jk} < m_{j_k}$  then
          begin  $m_{j_k} := m_{jk}$ ;  $w_{j_k} := i$ ;
            if  $k$  has not yet been reached then enter  $k$  at the end of the deque;
            if  $k$  has already been processed then enter  $k$  at the top of the deque
          end
        end
      end
    end
  end;

```

The great advantage of the above-mentioned technique is that errors in minimal distances are often corrected as soon as they are detected and not

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allowed to progress further. In practice this leads to a considerable reduction in CPU time, particularly for grid networks.

The additional bookkeeping work required costs less than the savings from the early processing of the nodes. This follows from a comparison of this technique with a more obvious strategy, where the deque is handled as a FIFO list (see [6]).

The network N is represented by its forward star structure (see [3]) which is also used as input to the subroutine. This input consists of three arrays: the array KF, containing information about where the data for each node begin in the other two arrays; the array FList, containing the numbers of the nodes which can be reached directly from each node; and the integer array DFList, containing the distances to each of these nodes.

The algorithm has been coded and tested in Algol 60, Algol-W, and Fortran IV for an IBM 370/158 (Berlin), an ICL 1907 (Braunschweig), and a CDC 6600 (Austin). It has been tested on different networks, real transportation/road networks and random road networks, constructed by network generators developed at Berlin and Austin. Independently, a variant of this algorithm has been tested at the Institute for Transportation Studies at the University of Leeds [8]. It has been compared with several versions of Dijkstra's algorithm [2] (with different binary tree structures for the nodes to be discussed), particularly Dial's algorithm [1]. The investigations have shown that the algorithm, based on d'Esopo's idea, is generally the most efficient in terms of CPU times (and core storage) for a large variety of networks. The code is highly insensitive to the magnitude of the distance measures.

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ALGORITHM

20						
5	2	4	3	5	6	
	8	8	8	9	14	
6	10	7	6	1	14	15
	9	9	10	8	14	14
4	6	1	11	8		
	8	8	8	10		
4	7	12	1	9		
	9	8	8	9		
5	1	9	8	13	18	
	9	9	8	9	14	
3	14	2	3			
	9	10	8			
3	16	2	4			
	8	9	9			
2	3	5				
	10	8				
5	4	5	18	1	12	
	9	9	9	14	14	
3	15	14	2			

		8	8	9	
2		3	19		
		8	20		
2		4	7		
		8	14		
2		5	17		
		9	9		
4	10	19		6	2
	8	10		9	14
2	10	2			
	8	14			
1		7			
	8				
2	13	20			
	9	8			
2	9	5			
	9	14			
2	14	11			
	10	20			
1	17				
	8				

- 1
- 1
- 2
- 3
- 4
- 5
- 6
- 7
- 8
- 9
- 10
- 11
- 12
- 13
- 14
- 15
- 16
- 17
- 18
- 19
- 20
- 0
- 3
- 20
- 15
- 10
- 5
- 6
- 9
- 0
- 11
- 2
- 0
- 20
- 16
- 8
- 4
- 2
- 1
- 0
- 0

```

      INTEGER FLIST(200),DFLIST(200),KF(51),MJ(50),WJ(50),W(50),H(50),
+      IN,INPUT,OUTP,TEMP(10)
C
C      THIS PROGRAM CALCULATES THE SHORTEST PATH LENGTHS FROM A SPECIFIC
C      NODE (J0) TO ALL OTHER NODES IN A NETWORK AND THE SHORTEST PATHS
C      BETWEEN THIS NODE (J0) AND OTHER NODES (KE).
C
C      FLIST(NUMBER OF EDGES)
C      DFLIST(NUMBER OF EDGES)
C      KF(NUMBER OF NODES+1)
C
C      FLIST, DFLIST, AND KF REPRESENT THE NETWORK

```

PAP00001
PAP00002
PAP00003
PAP00004
PAP00005
PAP00006
PAP00007
PAP00008
PAP00009
PAP00010
PAP00011
PAP00012

C		PAP00013
C	MJ(NUMBER OF NODES)	PAP00014
C	WJ(NUMBER OF NODES)	PAP00015
C	W(NUMBER OF NODES)	PAP00016
C	H(NUMBER OF NODES)	PAP00017
C		PAP00018
C	DATA INF,IN,INPUT,OUTP /1000000,4,5,6/	PAP00019
C		PAP00020
C	IN CHANNEL FOR INPUT J0	PAP00021
C	INPUT CHANNEL FOR INPUT NETWORK	PAP00022
C	OUTP CHANNEL FOR OUTPUT	PAP00023
C		PAP00024
C	INPUT OF THE NUMBER OF NODES N	PAP00025
C		PAP00026
1	READ(INPUT,100) N	PAP00027
	K=0	PAP00028
	KF(1)=0	PAP00029
C		PAP00030
C	INPUT OF THE NETWORK	PAP00031
C		PAP00032
	DO 4 I=1,N	PAP00033
	READ(INPUT,100) IT,(TEMP(J),J=1,IT)	PAP00034
	IF (IT.EQ.0) GO TO 3	PAP00035
	JA=K+1	PAP00036
	JB=K+IT	PAP00037
	JJ=1 .	PAP00038
	DO 2 J=JA,JB	PAP00039
	FLIST(J)=TEMP(JJ)	PAP00040
2	JJ=JJ+1	PAP00041
C		PAP00042
	READ(INPUT,100) IDUMMY,(DFLIST(J),J=JA,JB)	PAP00043
	K=K+IT	PAP00044
3	KF(I+1)=K	PAP00045
4	CONTINUE	PAP00046
C		PAP00047
C	INPUT OF START NODE J0	PAP00048
C		PAP00049
5	READ(IN,101) J0	PAP00050
	IF (J0.LE.0) STOP	PAP00051
C		PAP00052
	CALL SHPTHL(FLIST,DFLIST,KF,N,MJ,J0,WJ)	PAP00053
C		PAP00054
	WRITE(OUTP,200) J0	PAP00055
	IF (N.GT.0) WRITE(OUTP,201) (I,MJ(I),I=1,N)	PAP00056
C		PAP00057
C	INPUT OF END NODE KE	PAP00058
C		PAP00059
6	READ(IN,101) KE	PAP00060
	IF (KE.LE.0) GO TO 5	PAP00061
C		PAP00062
	CALL SHPATH(J0,KE,WJ,H,NI,W)	PAP00063
C		PAP00064
	WRITE(OUTP,202) J0,KE,MJ(KE)	PAP00065
	WRITE(OUTP,203) (W(I),I=1,NI)	PAP00066
	GO TO 6	PAP00067
C		PAP00068
100	FORMAT(10I8)	PAP00069
101	FORMAT(I4)	PAP00070
200	FORMAT(21H1DISTANCES FROM NODE ,I5/)	PAP00071
201	FORMAT(10(I7,1H-,I4))	PAP00072
202	FORMAT(/15H PATH FROM NODE,I7,9H TO NODE ,I7,13H (PATHLENGTH=,	PAP00073
	+ I7,1H))	PAP00074
203	FORMAT(1H ,20I6)	PAP00075
	END	PAP00076
	SUBROUTINE SHPTHL(FLIST,DFLIST,KF,N,MJ,J0,WJ)	SPL00001
	INTEGER FLIST(200),DFLIST(200),KF(51),MJ(50),NJ(50),WJ(50)	SPL00002
	DATA INF /1000000/	SPL00003
C		SPL00004
C	SHPTHL CALCULATES THE SHORTEST PATH LENGTHS (MJ) FROM A SPECIFIC	SPL00005
C	NODE (J0) TO ALL OTHER (N-1) NODES IN A NETWORK (FLIST,DFLIST,KF)	SPL00006
C	PREDECESSOR NODES ARE STORED IN WJ.	SPL00007
C		SPL00008
C	FLIST : FORWARD INDEX LIST	SPL00009
C	DFLIST: DISTANCE LIST	SPL00010

C	KF	: POINTER LIST FOR FLIST AND DFLIST	SPL00011
C	N	: NUMBER OF NODES	SPL00012
C	MJ	: ARRAY OF SHORTEST PATH LENGTHS	SPL00013
C	JO	: INITIAL NODE, FIRST NODE OF SHORTEST PATH	SPL00014
C	WJ	: ARRAY OF PREDECESSORS FOR SHORTEST PATH CONSTRUCTION	SPL00015
C	NJ	: DOUBLE ENDED QUEUE FOR NODE DISCUSSION	SPL00016
C	INF	: A LARGE NUMBER	SPL00017
C			SPL00018
C		INITIAL VALUES OF	SPL00019
C	FLIST,DFLIST,KF:	NETWORK FROM INPUT AND MAIN PROGRAM	SPL00020
C	N	: NUMBER OF NODES FROM INPUT AND MAIN PROGRAM	SPL00021
C	J0	: START NODE FROM INPUT AND MAIN PROGRAM	SPL00022
C			SPL00023
C		DO 1 I=1,N	SPL00024
C		MJ(I)=INF	SPL00025
C	1	NJ(I)=0	SPL00026
C		MJ(J0)=0	SPL00027
C			SPL00028
C	I	: INDEX FOR NODE DISCUSSION, NODE UNDER DISCUSSION	SPL00029
C	NT	: POINTER TO THE END OF DEQUE NJ	SPL00030
C	MJI	: LOCAL VARIABLE OF MJ(I)	SPL00031
C	KFI	: LOCAL VARIABLE OF KF(I)	SPL00032
C	KFI1	: LOCAL VARIABLE OF KF(I)+1	SPL00033
C	IR	: INDEX FOR ARRAY DISCUSSION	SPL00034
C	K	: SUCCESSOR OF NODE I	SPL00035
C	MJK	: LOCAL VARIABLE OF MJ(K)	SPL00036
C	NJI	: LOCAL VARIABLE OF NJ(I), THE NEXT NODE OF NJ TO BE TAKEN	SPL00037
C		UNDER DISCUSSION	SPL00038
C			SPL00039
C		NJ(J0)=INF	SPL00040
C		I=J0	SPL00041
C		NT=J0	SPL00042
C			SPL00043
C		OUTER LOOP	SPL00044
C		DISCUSSION OF NODES I	SPL00045
C			SPL00046
C	2	KFI=KF(I+1)	SPL00047
C		MJI=MJ(I)	SPL00048
C		KFI1=KF(I)+1	SPL00049
C			SPL00050
C		INNER LOOP	SPL00051
C		DISCUSSION OF SUCCESSORS K	SPL00052
C			SPL00053
C		IF (KFI1.GT.KFI) GO TO 6	SPL00054
C		DO 5 IR=KFI1,KFI	SPL00055
C		K=FLIST(IR)	SPL00056
C		MJK=MJI+DFLIST(IR)	SPL00057
C		NO DECREASE OF SHORTEST DISTANCES	SPL00058
C		IF (MJK.GE.MJ(K)) GO TO 5	SPL00059
C		DECREASE OF SHORTEST DISTANCES	SPL00060
C		MJ(K)=MJK	SPL00061
C		PREDECESSOR I OF NODE K	SPL00062
C		WJ(K)=I	SPL00063
C		NODE K ALREADY IN THE DEQUE NJ ?	SPL00064
C		IF (NJ(K)) 4,3,5	SPL00065
C		NODE K ADDED AT THE END OF THE DEQUE NJ	SPL00066
C	3	NJ(NT)=K	SPL00067
C		NT=K	SPL00068
C		NJ(K)=INF	SPL00069
C		GO TO 5	SPL00070
C		NODE K ADDED AT THE BEGINNING OF THE DEQUE NJ	SPL00071
C	4	NJ(K)=NJ(I)	SPL00072
C		NJ(I)=K	SPL00073
C	5	CONTINUE	SPL00074
C		NODE I TAKEN FROM THE BEGINNING OF THE DEQUE NJ	SPL00075
C	6	NJI=NJ(I)	SPL00076
C		NJ(I)=-NJI	SPL00077
C		I=NJI	SPL00078
C		IF (I.LT.INF) GOTO 2	SPL00079
C		RETURN	SPL00080
C		END	SPL00081
C			
C		SUBROUTINE SHPATH(J0,KE,WJ,H,NI,W)	PAT00001
C		INTEGER WJ(50),H(50),W(50)	PAT00002
C			PAT00003

C	SHPATH CALCULATES THE SHORTEST PATH BETWEEN THE TWO NODES J0 AND	PAT000004
C	KE. SHPATH USES THE INFORMATION IN WJ, THE LIST OF PREDECESSOR	PAT000005
C	NODES. THE NODES OF THE SHORTEST PATH ARE STORED IN W.	PAT000006
C		PAT000007
C	J0 : FIRST NODE OF THE SHORTEST PATH	PAT000008
C	KE : LAST NODE OF THE SHORTEST PATH	PAT000009
C	WJ : ARRAY OF PREDECESSORS FOR SHORTEST PATH CONSTRUCTION	PAT000010
C	H : AUXILIARY ARRAY FOR THE SHORTEST PATH	PAT000011
C	NI : NUMBER OF NODES OF THE SHORTEST PATH	PAT000012
C	W : THE SHORTEST PATH	PAT000013
C		PAT000014
C	I,J : LOCAL VARIABLE FOR NODE DISCUSSION	PAT000015
C		PAT000016
C	INITIAL VALUES OF	PAT000017
C	J0 : FIRST NODE FROM INPUT OR MAIN PROGRAM	PAT000018
C	KE : LAST NODE FROM INPUT OR MAIN PROGRAM	PAT000019
C	WJ : FROM SUBROUTINE SHPTH	PAT000020
C		PAT000021
	H(1)=KE	PAT000022
	I=1	PAT000023
	J=KE	PAT000024
	IF (J0.EQ.KE) GO TO 2	PAT000025
C		PAT000026
	1 I=I+1	PAT000027
	J=WJ(J)	PAT000028
	H(I)=J	PAT000029
	IF (J.NE.J0) GO TO 1	PAT000030
C		PAT000031
	2 NI=I	PAT000032
C		PAT000033
	DO 3 I=1,NI	PAT000034
	3 W(I)=H(NI+1-I)	PAT000035
C		PAT000036
	RETURN	PAT000037
	END	PAT000038

ALGORITHM 563

A Program for Linearly Constrained Discrete l_1 Problems

RICHARD H. BARTELS and ANDREW R. CONN
The University of Waterloo

Key Words and Phrases: numerical analysis, overdetermined linear systems, linear constraints, discrete l_1 approximation
CR Categories: 5.13, 5.41
Language: Fortran

DESCRIPTION

The subroutine CL1 given here is a complement of [1], where the theoretical development appears.

REFERENCES

1. BARTELS, R.H., AND CONN, A.R. Linearly constrained discrete l_1 problems. *ACM Trans. Math Softw.* 6, 4 (Dec. 1980), 594-608.

ALGORITHM

[Summary information and a part of the listing are printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

NAME(n): indicates a Fortran module with n records

NAME^T(n): indicates "NAME" is included for testing purposes

NAME^D(n): indicates "NAME" contains test data

Contents: MAIN^T(100), DATFL1^T(65), DATFL2^T(104), CL1(344),
SETUP(108), NEWPEN(57), UPDATE(52), MONIT(49),
FINDP(175), STEP(144), REFINE(48), DELCOL(51),
RESID(108), ADDCOL(70), OBJECT(72), GETV(91),
DKHEAP(108), UNIF01(88), ZDRCIN(164), ZDRCOU(166),
ZDRGIT(151), ZDRGNV(139), ZDRPOC(143), SASUM(69),
SAXPY(55), SCOPY(58), SDOT(54), SROTM(109),
SROTMG(181), SSCAL(41), DATAO^D(446)

	SUBROUTINE CL1 (NEQNS, NEQC, NIQC, NVAR, NACT, IFL, MXS, PSW,	CL102690
	* E, NER, X, F, ELIN, RES, INDX, W)	CL102700
C		CL102710
	INTEGER IFL, INDX(1), MXS, NACT, NEQC, NEQNS, NER, NIQC, NVAR	CL102720

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```

LOGICAL PSW
REAL E(NER,1),ELIN,F(1),RES(1),W(1),X(1)
C
C
C
C *****
C A PROGRAM FOR THE SOLUTION IN THE L1 SENSE
C OF A LINEAR-EQUATION SYSTEM
C (WITH OR WITHOUT LINEAR CONSTRAINTS).
C RICHARD H. BARTELS AND ANDREW R. CONN.
C LATEST UPDATE .... 13 APRIL, 1980.
C
C DEVELOPMENT OF THIS PROGRAM WAS SUPPORTED
C IN PART BY U. S. NSF GRANT DCR75-07817
C AND BY FUNDS FROM THE NATIONAL BUREAU OF STANDARDS,
C AND IN PART BY CANADIAN NRC GRANT A8639.
C
C +++++ PARAMETERS +++++
C -----
C          INPUT
C NAME  TYPE  SUBSCRIPT  OUTPUT  DESCRIPTION
C                               SCRATCH
C .....
C NEQNS INT.    NONE      IN      NUMBER OF EQUATIONS
C                               (MAY BE ZERO)
C
C NEQC  INT.    NONE      IN      NUMBER OF EQUALITY
C                               CONSTRAINTS
C                               (MAY BE ZERO)
C
C NIQC  INT.    NONE      IN      NUMBER OF INEQUALITY
C                               CONSTRAINTS
C                               (MAY BE ZERO)
C
C NVARs INT.    NONE      IN      NUMBER OF VARIABLES
C
C NACT  INT.    NONE      OUT     NUMBER OF ACTIVE
C                               EQUATIONS/CONSTRAINTS
C                               AT TERMINATION
C                               (IF ANY, THEIR ASSOCIATED
C                               COLUMN POSITIONS IN E WILL
C                               BE LISTED IN INDX(1)
C                               THROUGH INDX(NACT) )
C
C IFL   INT.    NONE      OUT     TERMINATION CODE
C                               (SEE BELOW)
C
C MXS   INT.    NONE      IN      MAXIMUM NUMBER OF STEPS
C                               ALLOWED
C
C PSW   LOGIC.  NONE      IN      PRINT SWITCH
C                               (SEE BELOW)
C
C E     REAL    2         IN      EQUATION/CONSTRAINT MATRIX
C                               THE FIRST NEQNS COLUMNS
C                               (SEE NOTE BELOW) SPECIFY
C                               EQUATIONS, THE REMAINING
C                               COLUMNS (IF ANY) SPECIFY
C                               CONSTRAINTS.
C
C NER   INT.    NONE      IN      ROW DIMENSION OF E
C
C X     REAL    1         IN      STARTING VALUES FOR THE
C                               UNKNOWNNS (USE ZEROS IF NO
C                               GUESS IS AVAILABLE)
C                               OUT     TERMINATION VALUES FOR
C                               THE UNKNOWNNS
C
C F     REAL    1         IN      EQUATION/CONSTRAINT
C                               RIGHT-HAND SIDES
C
C ELIN  REAL    NONE     OUT     L1 NORM OF EQUATION
C                               RESIDUALS AT TERMINATION
C
C RES   REAL    1         OUT     EQUATION/CONSTRAINT
C                               RESIDUALS AT TERMINATION

```

```

CL102730
CL102740
CL102750
CL102760
CL102770
CL102780
CL102790
CL102800
CL102810
CL102820
CL102830
CL102840
CL102850
CL102860
CL102870
CL102880
CL102890
CL102900
CL102910
CL102920
CL102930
CL102940
CL102950
CL102960
CL102970
CL102980
CL102990
CL103000
CL103010
CL103020
CL103030
CL103040
CL103050
CL103060
CL103070
CL103080
CL103090
CL103100
CL103110
CL103120
CL103130
CL103140
CL103150
CL103160
CL103170
CL103180
CL103190
CL103200
CL103210
CL103220
CL103230
CL103240
CL103250
CL103260
CL103270
CL103280
CL103290
CL103300
CL103310
CL103320
CL103330
CL103340
CL103350
CL103360
CL103370
CL103380
CL103390
CL103400
CL103410
CL103420
CL103430
CL103440
CL103450
CL103460
CL103470
CL103480

```

```

C      INDX INT.      1      OUT      INDEX VECTOR USED TO RECORD      CL103490
C      THE ORDER IN WHICH THE COLUMNS      CL103500
C      OF E ARE BEING PROCESSED      CL103510
C      CL103520
C      W      REAL      1      SCR.      WORKING STORAGE      CL103530
C      -----      CL103540
C      +++++ PURPOSE +++++      CL103550
C      -----      CL103560
C      THIS SUBROUTINE SOLVES THE      NEQNS BY NVARs      CL103570
C      SYSTEM OF EQUATIONS      CL103580
C      CL103590
C      (A-TRANSPPOSE) * X == B      CL103600
C      CL103610
C      SUBJECT TO THE      NEQC      CONSTRAINTS      CL103620
C      CL103630
C      (G-TRANSPPOSE) * X .EQ. H      CL103640
C      CL103650
C      AND THE      NIQC      INEQUALITY CONSTRAINTS      CL103660
C      CL103670
C      (C-TRANSPPOSE) * X .GE. D      CL103680
C      CL103690
C      FOR THE UNKNOWNs X(1),...X(NVARs).      CL103700
C      CL103710
C      THE PROBLEM MUST BE WELL-POSED, NONTRIVIAL      CL103720
C      AND OVERDETERMINED IN THE SENSE THAT      CL103730
C      CL103740
C      NVARs .GE. 1      CL103750
C      NEQNS .GE. 0      CL103760
C      NEQC .GE. 0      CL103770
C      NIQC .GE. 0      CL103780
C      NEQNS+NEQC+NIQC .GE. NVARs.      CL103790
C      CL103800
C      CL103810
C      FURTHER, NO COLUMN OF A, G OR C SHOULD BE ZERO.      CL103820
C      IF THESE CONDITIONS ARE NOT MET, THE PROGRAM      CL103830
C      WILL TERMINATE WITHOUT PERFORMING ANY SUBSTANTIVE      CL103840
C      COMPUTATIONS.      CL103850
C      CL103860
C      A POINT X IS A SOLUTION IF IT MINIMIZES THE EQUATION      CL103870
C      RESIDUALS FROM AMONG ALL POINTS WHICH SATISFY THE      CL103880
C      CONSTRAINTS. AT ANY (NONDEGENERATE) SOLUTION      CL103890
C      THERE WILL BE NACT EQUATIONS AND CONSTRAINTS      CL103900
C      WHOSE RESIDUALS      CL103910
C      CL103920
C      (A(I)-TRANSPPOSE) * X - B(I)      CL103930
C      CL103940
C      (G(I)-TRANSPPOSE) * X - H(I)      CL103950
C      CL103960
C      AND      CL103970
C      CL103980
C      (C(I)-TRANSPPOSE) * X - D(I)      CL103990
C      CL104000
C      ARE ZERO.      CL104010
C      CL104020
C      THE COLUMNS OF (A,G,C) CORRESPONDING TO THE ZERO RESIDUALS      CL104030
C      ARE REFERRED TO AS ACTIVE COLUMNS THROUGHOUT THIS LISTING.      CL104040
C      THE NUMBERS OF THE ACTIVE COLUMNS ARE MAINTAINED AS THE      CL104050
C      ENTRIES 1,...,NACT OF THE ARRAY INDX.      CL104060
C      CL104070
C      A SOLUTION X IS FOUND BY MINIMIZING A PIECEWISE      CL104080
C      LINEAR PENALTY FUNCTION FORMED FROM THE L1      CL104090
C      NORM OF THE EQUATION RESIDUALS AND THE SUM OF THE      CL104100
C      INFEASIBILITIES IN THE CONSTRAINTS.      CL104110
C      THE MINIMIZATION PROCEEDS IN A STEP-BY-STEP      CL104120
C      FASHION, TERMINATING AFTER A FINITE NUMBER OF STEPS.      CL104130
C      CL104140
C      NOTE THAT A, G AND C APPEAR TRANSPPOSED IN THE      CL104150
C      PROBLEM FORMULATION. HENCE IT IS THE COLUMNS OF (A,G,C)      CL104160
C      WHICH DEFINE THE EQUATIONS AND CONSTRAINTS RESPECTIVELY.      CL104170
C      CL104180
C      THE ARRAY E IS A COMPOSITE OF A, G AND C      CL104190
C      AND F IS A COMPOSITE OF B, H AND D.      CL104200
C      E SHOULD CONTAIN A AS ITS FIRST NEQNS COLUMNS.      CL104210
C      IT SHOULD CONTAIN G AS ITS NEXT NEQC COLUMNS AND      CL104220
C      CONTAIN C AS ITS REMAINING NIQC COLUMNS.      CL104230
C      SIMILARLY F SHOULD CONTAIN B AS ITS FIRST      CL104240

```

```

C      NEQNS COMPONENTS, H AS ITS NEXT NEQC COMPONENTS          CL104250
C      AND D AS ITS LAST NIQC COMPONENTS.                       CL104260
C      -----                                                  CL104270
C      +++++ ARRAYS +++++                                       CL104280
C      -----                                                  CL104290
C      E IS TO BE DIMENSIONED AT LEAST N BY M,                  CL104300
C      X                      AT LEAST N,                       CL104310
C      F                      AT LEAST M,                       CL104320
C      RES                   AT LEAST M,                       CL104330
C      INDX                  AT LEAST M,                       CL104340
C      W                      AT LEAST ((3*N*N+11*M+2)/2) + (2*M). CL104350
C                                                                CL104360
C                                                                CL104370
C                                                                CL104380
C                                                                CL104390
C                                                                CL104400
C                                                                CL104410
C                                                                CL104420
C      +++++ INITIALIZATION +++++                               CL104430
C      -----                                                  CL104440
C      THE USER MUST INITIALIZE                                CL104450
C                                                                CL104460
C      NEQNS,NEQC,NIQC,NVARS,MXS,PSW,E,NER,X,F .              CL104470
C                                                                CL104480
C      THE FOLLOWING ARE SET BY CL1                             CL104490
C      AND DO NOT REQUIRE INITIALIZATION                      CL104500
C                                                                CL104510
C      NACT,INDX,RES .                                         CL104520
C                                                                CL104530
C      THE ARRAY W IS USED AS SCRATCH SPACE.                  CL104540
C      -----                                                  CL104550
C      +++++ TERMINATION CODES AND INTERMEDIATE PRINTING +++++ CL104560
C      -----                                                  CL104570
C      MXS SETS A LIMIT ON THE NUMBER OF MINIMIZATION STEPS TO BE CL104580
C      TAKEN.                                                  CL104590
C                                                                CL104600
C      UPON TERMINATION IFL WILL BE SET ACCORDING TO          CL104610
C      THE FOLLOWING CODE ...                                   CL104620
C                                                                CL104630
C      IFL = 1 ... SUCCESSFUL TERMINATION.                     CL104640
C                                                                CL104650
C      IFL = 2 ... UNSUCCESSFUL TERMINATION.                   CL104660
C      CONSTRAINTS CANNOT BE SATISFIED.                       CL104670
C      PROBLEM IS INFEASIBLE.                                 CL104680
C                                                                CL104690
C      IFL = 3 ... LIMIT IMPOSED BY MXS REACHED               CL104700
C      WITHOUT FINDING A SOLUTION.                             CL104710
C                                                                CL104720
C      IFL = 4 ... PROGRAM ABORTED.                            CL104730
C      NUMERICAL DIFFICULTIES                                 CL104740
C      DUE TO ILL-CONDITIONING.                               CL104750
C                                                                CL104760
C      IFL = 5 ... NEQNS, NVARS, NEQC AND/OR                  CL104770
C      NIQC HAVE IMPROPER VALUES                             CL104780
C      OR E CONTAINS A ZERO COLUMN.                            CL104790
C                                                                CL104800
C      IN ALL CASES THE OUTPUT PARAMETERS X,ELIN AND RES       CL104810
C      WILL CONTAIN THE VALUES WHICH THEY REACHED AT TERMINATION. CL104820
C                                                                CL104830
C      INTERMEDIATE PRINTING WILL BE TURNED OFF IF PSW = .FALSE. CL104840
C      ON THE OTHER HAND, DETAILS OF EACH MINIMIZATION CYCLE CL104850
C      WILL BE PRINTED IF PSW IS SET TO .TRUE.                CL104860
C      -----                                                  CL104870
C                                                                CL104880
C      +++++ REMARKS AND USER CAUTIONS +++++                   CL104890
C      -----                                                  CL104900
C      1. BEYOND SOME PRECAUTIONARY STEPS TAKEN IN            CL104910
C      CERTAIN DIVISIONS, NO SPECIAL                           CL104920
C      OVERFLOW/UNDERFLOW PROTECTION IS PROVIDED.            CL104930
C      2. ALL TOLERANCES FOR CHECKING ZEROS AND LINEAR        CL104940
C      DEPENDENCIES ARE DETERMINED FROM THE QUANTITY          CL104950
C      EPS WHICH APPEARS IN DATA DECLARATIONS IN CL1         CL104960
C      AND SEVERAL OF ITS SUBROUTINES. EPS CAN BE SET TO THE CL104970
C      LEAST POSITIVE NUMBER SATISFYING (1.0 + EPS) .GT. 1.0 CL104980
C      IN THE PRECISION OF ARITHMETIC BEING USED. WITH THIS CL104990
C      SETTING, CL1 USES AN EXTREMELY STRICT ZERO TOLERANCE. CL105000

```

C	FOR A MORE FORGIVING VERSION OF CL1, EPS MAY BE REMOVED	CL105010
C	FROM THE DATA DEFINITIONS AND INCLUDED AS A USER-SPECIFIED	CL105020
C	ZERO-TESTING PARAMETER IN THE ARGUMENT LIST. IN SUCH AN	CL105030
C	EVENT, IF THE PROBLEM DATA IS GIVEN TO NDIG SIGNIFICANT	CL105040
C	DIGITS, THEN $10.0^{**(-NDIG)}$ IS A REASONABLE CHOICE	CL105050
C	FOR THE VALUE OF EPS.	CL105060
C	OVERFLOW CHECKING PRIOR TO DIVISION IS	CL105070
C	DONE USING THE QUANTITY BIG, ALSO SPECIFIED AS DATA.	CL105080
C	BIG SHOULD BE THE LARGEST REPRESENTABLE FLOATING	CL105090
C	POINT NUMBER.	CL105100
C	3. THIS IS A SINGLE PRECISION VERSION OF CL1.	CL105110
C	TO CHANGE THIS CODE INTO DOUBLE PRECISION ...	CL105120
C	A. CHANGE ALL OCCURRENCES OF - REAL -	CL105130
C	DECLARATIONS TO - DOUBLE PRECISION -	CL105140
C	B. CHANGE ALL OCCURRENCES OF - SYSTEM ROUTINES -	CL105150
C	(AS LISTED IN THE HEADING OF EACH SUBROUTINE)	CL105160
C	TO THEIR CORRESPONDING DOUBLE PRECISION VERSIONS	CL105170
C	C. CHANGE ALL OCCURRENCES OF THE STRINGS E+ AND E-	CL105180
C	TO D+ AND D- RESPECTIVELY	CL105190
C	D. CHANGE ALL BASIC LINEAR ALGEBRA ROUTINES (BLAS)	CL105200
C	TO THEIR DOUBLE-PRECISION EQUIVALENTS	CL105210
C	E. BOTH EPS AND BIG WILL HAVE TO BE CHANGED	CL105220
C	TO THEIR DOUBLE PRECISION EQUIVALENTS	CL105230
C	F. THE REFERENCES TO - IFIX - IN SUBROUTINES	CL105240
C	- RESID - AND - GETV - MUST BE CHANGED	CL105250
C	FROM THE FORM	CL105260
C	IFIX(FLOAT(K)*UNIF(...))	CL105270
C	TO THE FORM	CL105280
C	IFIX(FLOAT(K)*SNGL(UNIF(...)))	CL105290
C	G. THE REFERENCE TO - FLOAT - IN SUBROUTINE	CL105300
C	- RESID - MUST BE CHANGED FROM THE FORM	CL105310
C	FLOAT(...)	CL105320
C	TO THE FORM	CL105330
C	DBLE(FLOAT(...))	CL105340
C	H. REMOVE THESE COMMENT CARDS (3., 3.A.-3.H.).	CL105350
C	-----	CL105360

Editor's Note

Owing to the increasing length of algorithms accepted for publication in ACM Transactions, it has become impractical to supply lengths of individual procedures. Instead, lengths of entire algorithms are given on the Algorithms order form.

ALGORITHM 564

A Test Problem Generator for Discrete Linear L_1 Approximation Problems

K.L. HOFFMAN and D.R. SHIER
National Bureau of Standards

Key Words and Phrases: L_1 approximation, least absolute deviation, problem generator, test data
CR Categories: 5.13, 5.41, 5.5
Language: Fortran

DESCRIPTION

The algorithm given here is a complement to [1], where its theoretical development and implementation are described.

REFERENCES

1. HOFFMAN, K.L., AND SHIER, D.R. A test problem generator for discrete linear L_1 approximation problems. *ACM Trans. Math. Softw.* 6, 4 (Dec. 1980), 587-593.

ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```
C DESCRIPTION:
C
C   THIS SUBROUTINE GENERATES DATA SETS WHICH CAN BE
C   USED FOR TESTING L1-APPROXIMATION (LEAST ABSOLUTE
C   DEVIATION CURVE-FITTING) COMPUTER CODES.  THE USER
C   CAN SPECIFY
C
C   * PROBLEM SIZE
C   * SOLUTION VECTOR
C   * STATISTICAL DISTRIBUTION FOR COLUMN ELEMENTS
C   * ROW REPETITIONS
C   * DEGENERACY
C   * RANK LOSS
C   * STATISTICAL DISTRIBUTION FOR RESIDUALS
C
```

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```

C   ON OUTPUT, X = XOPT IS AN OPTIMAL SOLUTION TO THE
C   GENERATED L1-APPROXIMATION PROBLEM:
C
C
C   MINIMIZE // X*A - RHS // ,
C
C
C   WHERE //...// IS THE L1-NORM AND THE TRANSPOSE OF
C   A REPRESENTS THE DESIGN MATRIX.  FURTHERMORE,
C   X = XOPT IS GUARANTEED TO BE THE UNIQUE OPTIMAL
C   SOLUTION IF RLOSS=0.
C
C INPUT ARGUMENTS:
C
C   NOBS -- NUMBER OF OBSERVATIONS
C   NPAR -- NUMBER OF PARAMETERS
C   NXGEN -- INDICATES WHETHER SOLUTION X IS SPECIFIED
C           AS INPUT OR GENERATED
C           = 0 IF SPECIFIED AS INPUT
C           = 1 IF RANDOMLY GENERATED FROM NORMAL
C           = 2 IF RANDOMLY GENERATED FROM UNIFORM
C   XOPT -- REAL ARRAY OF DIMENSION NPAR - RLOSS WHICH
C           SPECIFIES THE OPTIMAL SOLUTION IF NXGEN IS
C           SET EQUAL TO ZERO
C   XMEAN -- MEAN OF NORMAL DISTRIBUTION IF NXGEN=1,
C           LOWER LIMIT OF UNIFORM DISTRIBUTION IF
C           NXGEN=2
C   XVAR -- STANDARD DEVIATION OF NORMAL DISTRIBUTION
C           IF NXGEN=1, UPPER LIMIT OF UNIFORM DISTRI-
C           BUTION IF NXGEN=2
C   NDIST -- ARRAY OF DIMENSION NPAR - RLOSS WHICH
C           SPECIFIES THE STATISTICAL DISTRIBUTION
C           FOR EACH COLUMN OF THE DESIGN MATRIX
C           = 0 IF NORMAL
C           = 1 IF UNIFORM
C           = 2 IF ZERO-ONE
C   CLMEAN -- ARRAY OF DIMENSION NPAR - RLOSS WHICH
C           SPECIFIES FOR EACH COLUMN THE MEAN (IF
C           NDIST=0), LOWER LIMIT (IF NDIST=1), OR
C           PROPORTION P OF ZERO ENTRIES (IF NDIST=2)
C   CLMVAR -- ARRAY OF DIMENSION NPAR - RLOSS WHICH
C           SPECIFIES FOR EACH COLUMN THE STANDARD
C           DEVIATION (IF NDIST=0), THE UPPER LIMIT
C           (IF NDIST=1); NOT USED IF NDIST=2
C
C   ** NOTE **   A CONSTANT COLUMN OF 1'S CAN BE GENERATED
C               USING CLMEAN = CLMVAR = 1.0 AND NDIST = 1,
C               FOR EXAMPLE.
C
C   NREPS -- NUMBER OF TIMES EACH UNIQUE ROW IS TO
C           APPEAR IN THE GENERATED DESIGN MATRIX
C   NDEG -- NUMBER OF ZERO RESIDUALS OUTSIDE OF THE
C           BASIS
C   RLOSS -- NUMBER OF DEPENDENT COLUMNS ADDED
C   NIND -- VALUE SUCH THAT THE DEPENDENT COLUMNS ARE
C           GENERATED BY SUMMING NIND SUCCESSIVE
C           COLUMNS
C   YDIST -- INTEGER WHICH SPECIFIES THE DISTRIBUTION
C           OF (NONZERO) RESIDUALS ASSOCIATED WITH
C           THE NONACTIVE CONSTRAINTS
C           = 0 IF NORMAL
C           = 1 IF UNIFORM
C   YMEAN -- MEAN OF RESIDUAL DISTRIBUTION (IF YDIST=0),
C           OR LOWER LIMIT (IF YDIST=1)
C   YVAR -- STANDARD DEVIATION OF RESIDUAL DISTRIBUTION
C           (IF YDIST=0), OR UPPER LIMIT (IF YDIST=1)
C   ISEED -- THE RANDOM NUMBER SEED USED TO
C           INITIATE THE RANDOM NUMBER GENERATOR
C
C   ** NOTE **   IF NXGEN, NDIST OR YDIST ARE OUTSIDE THE

```

```
C          INDICATED RANGES, THE REQUIRED DISTRIBUTION
C          IS DEFAULTED TO THAT OF A NORMAL.
C
C OUTPUT ARGUMENTS:
C
C   XOPT -- OPTIMAL SOLUTION VECTOR OF DIMENSION NPAR
C   A -- GENERATED TRANSPOSE OF THE DESIGN MATRIX, WITH
C        NPAR ROWS AND NOBS COLUMNS
C   RHS -- RIGHT HAND SIDE (OR Y) VECTOR OF DIMENSION
C        NOBS
C   IACT -- VECTOR OF DIMENSION NPAR - RLOSS WHICH
C          CONTAINS INDICES OF ACTIVE CONSTRAINT ROWS
C   SUMRES -- OBJECTIVE FUNCTION VALUE, OR SUM OF
C            ABSOLUTE VALUES OF RESIDUALS
C   ISEED --- THE RANDOM NUMBER SEED AVAILABLE
C            UPON TERMINATION OF GENERATION PROCESS
C   IERR -- ERROR FLAG, UPON RETURN
C           = 0 NORMAL EXECUTION
C           = 1 FATAL ERROR
C
C RESTRICTIONS:
C
C   1 .LE. NOBS .LE. 400
C   1 .LE. NPAR .LE. 25
C   1 .LE. NREPS .LE. NOBS
C   0 .LE. NDEG .LE. NOBS-NPAR+RLOSS-2
C   0 .LE. RLOSS .LE. NPAR-1
C   0 .LE. NIND .LE. NPAR-RLOSS
C   NOBS+RLOSS-NPAR-NDEG MUST BE EVEN
C   NREPS MUST DIVIDE NOBS
```


ALGORITHM 565

PDETWO/PSETM/GEARB: Solution of Systems of Two-Dimensional Nonlinear Partial Differential Equations [D3]

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and

RICHARD F. SINCOVEC

Boeing Computer Services Company

Key Words and Phrases: partial differential equations, method of lines, finite differences, ordinary differential equations

CR Categories: 5.17

Language: FORTRAN

DESCRIPTION

The algorithm presented here combines PDETWO, PSETM[1], and a modified version of GEARB[2] to form a complete partial differential equations package for systems of time-dependent nonlinear partial differential equations defined over a two-dimensional rectangular region. The descriptions of PDETWO and PSETM, test results, references, and the use of this algorithm are contained in the authors' paper [2]. The ordinary differential equations package, GEARB, is described by Hindmarsh [1], so we only comment on those modifications that we made to GEARB for this algorithm. The principal modification is the addition of a new subroutine to set pointers for dynamic dimensioning of arrays and to check the validity of some of the user input parameters. The interpolation routine in GEARB was also modified so that it would return correctly interpolated solution values at the user-specified output time for time-dependent boundary conditions.

Machine-readable documentation within the package serves as a complete user's guide. The comments in subroutine PDETWO give the essential details on the use of this package, including descriptions of the user-required subroutines. Additional user information is contained in subroutine DRIVEP, the driver routine for the package.

REFERENCES

1. HINDMARSH, H.C. GEARB: Solution of ordinary differential equations having banded Jacobian. Rep. UCID-30059, Rev. 2, Lawrence Livermore Lab., Livermore, Calif., June 1977.

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2. MELGAARD, D.K., AND SINCOVEC, R.F. General software for two-dimensional nonlinear partial differential equations. *ACM Trans. Math. Softw.* 7, 1 (March 1981), 106-125.

ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service].

```

SUBROUTINE PDETWO (NPDE, NX, NY, X, Y, T, U, DUDT, DVS, DVST, DV, DH,
*             AH, BH, CH, AV, BV, CV, UX, UY, DXI, DXIR, DXIC,
*             XAVG, UDVA, UDHR, UDV, UDHL, UAVGH, UAVGV)
C
C*****
C
C      PDETWO IS AN INTERFACE DESIGNED TO SOLVE A SYSTEM OF SECOND
C ORDER TIME DEPENDENT PARTIAL DIFFERENTIAL EQUATIONS (PDE*S)
C DEFINED OVER A RECTANGLE IN CONJUNCTION WITH AN ORDINARY
C DIFFERENTIAL EQUATION (ODE) INTEGRATOR PACKAGE (PRIMARILY A MODI-
C FIED VERSION OF GEARE(1)). IT IS AN EXTENSION OF PDEONE (2,3), AN
C INTERFACE DEVELOPED FOR THE SOLUTION OF ONE-DIMENSIONAL SYSTEMS OF
C PDE*S. THIS EXTENSION CONSISTS OF CHANGING FROM THREE POINT
C DIFFERENCING TO FIVE POINT DIFFERENCING IN ORDER TO APPROXIMATE
C THE SPATIAL DERIVATIVES IN THE TWO DIMENSIONAL PDE SYSTEM. TO
C SOLVE A PDE SYSTEM USING THIS SOFTWARE INTERFACE IN CONJUNCTION
C WITH AN ODE INTEGRATOR, THE USER IS REQUIRED ONLY TO PROVIDE
C THE SPATIAL MESH AND DEFINE THE PROBLEM TO BE EVALUATED. THE
C INTERFACE WILL THEN FORM AND EVALUATE A SEMIDISCRETE APPROXI-
C MATION OF THIS SYSTEM OF PDE*S AND THUS PERMIT ONE TO SOLVE
C THE ORIGINAL PDE SYSTEM AS A SYSTEM OF TIME DEPENDENT ODE*S
C USING AN ODE INTEGRATOR.
C
C      TO CREATE A DOUBLE PRECISION VERSION OF PDETWO..
C CHANGE THE REAL STATEMENT BELOW.
C
C*****
C PROBLEM DEFINITION
C*****
C
C      PDETWO IS DESIGNED TO BE USED TO SOLVE A COUPLED SYSTEM OF
C NPDE PDE*S. THE L-TH PDE OF THIS SYSTEM IS OF THE FORM ..
C
C      DUDT(L) = F ( T, X, Y, U, UX, UY, DUXX, DUY )
C
C WHERE
C
C      DUDT(L) REPRESENTS THE FIRST PARTIAL OF THE L-TH COMPONENT
C OF U WITH RESPECT TO T,
C T IS THE CURRENT TIME,
C X,Y DEFINE THE POSITION IN THE HORIZONTAL AND VERTICAL
C DIRECTIONS RESPECTIVELY,
C
C      U = ( U(1), ... , U(NPDE) ),
C      UX = ( UX(1), ... , UX(NPDE) ),
C      UY = ( UY(1), ... , UY(NPDE) ),
C
C      DUXX IS AN NPDE BY NPDE ARRAY SUCH THAT
C
C          D
C      DUXX(L,K) = -- (DH(L,K)*UX(K)),
C          DX
C      THE L-TH ROW OF DUXX CORRESPONDS TO THE L-TH PDE,
C
C      DUY IS AN NPDE BY NPDE ARRAY SUCH THAT
C
C          D
C      DUY(L,K) = -- (DV(L,K)*UY(K)),
C          DY
C      THE L-TH ROW OF DUY CORRESPONDS TO THE L-TH PDE,
C
C
C

```

```

C   AND
C
C   U(K)   IS THE VALUE OF THE SOLUTION FOR THE K-TH PDE AT
C           (T, X, Y),
C   UX(K)  IS THE FIRST PARTIAL DERIVATIVE OF U(K) WITH RESPECT
C           TO X,
C   UY(K)  IS THE FIRST PARTIAL DERIVATIVE OF U(K) WITH RESPECT
C           TO Y,
C   DH(L,K) IS THE DIFFUSION COEFFICIENT IN THE HORIZONTAL
C           DIRECTION FOR THE L-TH PDE ASSOCIATED WITH U(K),
C   DV(L,K) IS THE DIFFUSION COEFFICIENT IN THE VERTICAL
C           DIRECTION FOR THE L-TH PDE ASSOCIATED WITH U(K).
C
C   DH(L,K) AND DV(L,K) MAY BE FUNCTIONS OF T, X, Y AND U.
C
C   HORIZONTAL BOUNDARY CONDITIONS
C
C   AH(L)*U(L) + BH(L)*UY(L) = CH(L)
C
C   VERTICAL BOUNDARY CONDITIONS
C
C   AV(L)*U(L) + BV(L)*UX(L) = CV(L)
C
C   AH(L), BH(L) AND CH(L) (AV(L), BV(L) AND CV(L)) ARE AT LEAST
C   PIECEWISE CONTINUOUS FUNCTIONS OF THEIR RESPECTIVE VARIABLES.
C   IF BH(L) .NE. 0 (BV(L) .NE. 0) THEN AH(L), BH(L) AND CH(L)
C   (AV(L), BV(L) AND CV(L)) MAY BE FUNCTIONS OF T, X, Y AND U
C   BUT OTHERWISE THEY MAY ONLY BE FUNCTIONS OF T, X AND Y. NULL
C   BOUNDARY CONDITIONS ARE NOT ALLOWED SINCE PDETWO WAS NOT
C   DESIGNED TO SOLVE HYPERBOLIC PDE*S. A ZERO DIVIDE WILL OCCUR
C   IF A NULL BOUNDARY CONDITION IS SPECIFIED.
C
C   INITIAL CONDITIONS
C
C   THE INITIAL SOLUTION IS DEFINED FOR EACH U(K) TO BE A
C   KNOWN FUNCTION OF X AND Y FOR THE INITIAL TIME T = T0.
C   THE INITIAL CONDITIONS NEED NOT BE CONSISTENT WITH THE
C   BOUNDARY CONDITIONS AND MAY CONTAIN DISCONTINUITIES.
C
C*****
C   USER SUPPLIED ROUTINES
C*****
C
C   THE USER MUST PROVIDE A MAIN PROGRAM AND FIVE SUBROUTINES
C   BNDRYH, BNDRYV, DIFFV, DIFFH AND F. THESE ROUTINES ARE ALL THAT
C   IS REQUIRED TO DEFINE COMPLETELY THE PROBLEM GIVEN ABOVE.
C
C   1) THE MAIN PROGRAM DEFINES THE SPATIAL MESH, THE INITIAL
C   CONDITIONS AND THE PARAMETERS FOR THE ODE INTEGRATOR, CALLS
C   THE INTEGRATOR AND PRINTS OR PLOTS THE RESULTS. THE MAIN
C   PROGRAM SHOULD BE CONSTRUCTED AS FOLLOWS ..
C
C   DIMENSION U(***,*,**),X(*),Y(**)
C   DIMENSION WORK(*****),IWORK(****)
C
C   FOR THE DIMENSIONS ABOVE ENTER THE ACTUAL NUMERICAL
C   VALUES FOR * = NX, ** = NY, *** = NPDE, **** = NODE, AND
C   ***** = NPDE * ( NPDE*(NX*2+3) + 13 + NX ) + NX*4 + 4*NODE
C   + LPW + LU
C   WHERE
C   NODE = NPDE*NX*NY
C   LPW  = 1                               IF MITER = 0
C        = (3*(NX+1)*NPDE-2)*NODE        IF MITER = 1,2
C        = NODE                             IF MITER = 3
C   LU  = NODE*(MORDER+1)
C   SEE MF AND MORDER BELOW FOR THE DEFINITION OF MITER AND
C   MORDER.
C
C   DEFINE ..
C
C   1) THE NUMBER OF MESH POINTS IN THE X (HORIZONTAL) DIRECTION,
C   NX .GE. 3 AND THE NUMBER OF MESH POINTS IN THE Y (VERTICAL)
C   DIRECTION, NY .GE. 3 ( TO CONSERVE STORAGE ORIENT THE
C   PROBLEM SO THAT NX .LE. NY ),
C   2) THE MESH POINTS (I.E. X(1) .LT. X(2) .LT. ... .LT. X(NX)

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```

C      AND Y(1) .LT. Y(2) .LT. ... .LT. Y(NY)),
C      3) NPDE, THE NUMBER OF PARTIAL DIFFERENTIAL EQUATIONS,
C      4) INITIAL VALUES FOR ALL OF U,
C      5) MORDER, THE MAXIMUM ORDER OF THE METHOD USED IN THE
C      MODIFIED GEARB. MORDER MUST BE LESS THAN OR EQUAL TO 12
C      IF METH = 1 OR IF METH = 2, IT MUST BE LESS THAN OR EQUAL
C      TO 5. SEE MF BELOW FOR THE DEFINITION OF METH.
C      6) THE PARAMETERS FOR THE CALL TO THE INTEGRATOR, AND
C      7) THE FIRST SIX LOCATIONS OF THE ARRAY IWORK AS FOLLOWS.
C      IWORK(1) = NPDE
C      IWORK(2) = NX
C      IWORK(3) = NY
C      IWORK(4) = MORDER
C      IWORK(5) = NRWK
C      IWORK(6) = NRIWK
C      WHERE NRWK AND NRIWK ARE THE LENGTHS OF THE ARRAYS WORK
C      AND IWORK RESPECTIVELY. THEY SHOULD BE AT LEAST EQUAL
C      TO ***** AND **** AS DEFINED IN THE DESCRIPTION OF THE
C      DIMENSIONS.
C
C      IN DETERMINING THE MESH SPACING, THE USER SHOULD BE AWARE
C      THAT THE DIFFERENCE APPROXIMATIONS AT THE INTERIOR POINTS
C      ARE SECOND ORDER ONLY FOR UNIFORM MESHES. FOR NONUNIFORM
C      MESHES, ONLY FIRST ORDER APPROXIMATIONS SHOULD BE EXPECTED.
C      IN UNUSUAL SITUATIONS (SEE REMARK IN SECTION 3 OF THE
C      ACCOMPANYING PAPER) INVOLVING COUPLED SYSTEMS OF EQUATIONS,
C      CERTAIN APPROXIMATIONS AT THE BOUNDARY ARE ONLY FIRST ORDER.
C      IN THESE UNUSUAL SITUATIONS THE ERROR CAN BE MINIMIZED BY
C      CHOOSING SMALL MESH SPACINGS NEXT TO THE BOUNDARY. IN ANY
C      CASE THE USER IS ADVISED TO CHOOSE A MESH WHICH IS LOCALLY
C      NEARLY UNIFORM.
C
C      FOR THE MODIFIED GEARB THE PARAMETERS INCLUDE THE DESIRED
C      OUTPUT TIME (TOUT), THE DESIRED LOCAL ACCURACY (EPS), THE
C      NUMBER OF ODE*S (NODE = NPDE*NX*NY), THE INITIAL TIME (T0),
C      THE INITIAL TIME STEP SIZE (H), THE TYPE OF CALL BEING MADE
C      (INDEX) AND THE TYPE OF INTEGRATION METHOD DESIRED (MF).
C      MF HAS TWO DECIMAL DIGITS, METH AND MITER (MF = 10*METH +
C      MITER). METH IS THE BASIC METHOD INDICATOR..
C      METH = 1 MEANS THE ADAMS METHODS.
C      METH = 2 MEANS THE BACKWARD DIFFERENTIATION FORMULAS
C      (BDF), OR STIFF METHODS OF GEAR.
C      MITER IS THE ITERATION METHOD INDICATOR..
C      MITER = 0 MEANS FUNCTIONAL ITERATION (NO PARTIAL
C      DERIVATIVES NEEDED).
C      MITER = 1 MEANS THE CHORD METHOD WITH AN ANALYTIC
C      JACOBIAN SUPPLIED IN THE USER DEFINED
C      ROUTINE PDB. THIS METHOD IS IN GEARB,
C      BUT IT SHOULD BE AVOIDED.
C      MITER = 2 MEANS THE CHORD METHOD WITH THE JACO-
C      BIAN CALCULATED IN PSETM. THIS IS THE
C      ONLY METHOD WHERE PSETM IS USEFUL.
C      MITER = 3 MEANS THE CHORD METHOD WITH THE JACO-
C      BIAN REPLACED BY A DIAGONAL APPROXI-
C      MATION BASED ON A DIRECTIONAL DERIV-
C      ATIVE. THIS METHOD IS IN GEARB, BUT
C      IS NOT GENERALLY USEFUL IN SOLVING PDE*S.
C      SEE COMMENTS IN DRIVEP FOR ADDITIONAL INFORMATION ON THESE
C      PARAMETERS. NATURALLY THESE PARAMETERS ARE DEPENDENT ON THE
C      INTEGRATOR BEING USED. FINALLY THE USER SHOULD CALL THE
C      INTEGRATOR.
C
C      CALL DRIVEP (NODE,T0,H,U,TOUT,EPS,MF,INDEX,WORK,IWORK,X,Y)
C
C      THE INTEGRATOR WILL RETURN THE SOLUTIONS (U) AT T = TOUT
C      TO THE MAIN PROGRAM TO BE PRINTED OR PLOTTED. IF A CONTINUA-
C      TION TO ANOTHER TOUT IS DESIRED, SIMPLY RESET TOUT AND CALL
C      DRIVEP AGAIN.
C
C      STOP
C      END

```



```

C
C 2) THE BOUNDARY SUBROUTINES, BNDRYH AND BNDRYV PROVIDE
C PDETWO WITH THE COEFFICIENTS FOR THE HORIZONTAL AND
C VERTICAL BOUNDARY CONDITIONS RESPECTIVELY. THE CONSTRUCTION
C OF BNDRYV IS ANALOGOUS TO THE CONSTRUCTION OF BNDRYH. THE
C USER SHOULD CONSTRUCT BNDRYH AS FOLLOWS ..
C
C SUBROUTINE BNDRYH (T,X,Y,U,AH,BH,CH,NPDE)
C DIMENSION U(NPDE), AH(NPDE), BH(NPDE), CH(NPDE)
C
C THE INCOMING PARAMETERS X AND Y REPRESENT ANY POINT
C OF X(J) (J = 1,2,...,NX) AND EITHER Y(1) OR Y(NY) RESPECTIVELY.
C DEFINE THE FUNCTIONS AH(K), BH(K) AND CH(K) (K = 1,2,...,NPDE)
C FOR THE LOWER (Y=Y(1)) AND UPPER (Y=Y(NY)) BOUNDARIES. NULL
C BOUNDARY CONDITIONS ARE NOT ALLOWED AND IF SPECIFIED WILL
C RESULT IN A ZERO DIVIDE.
C
C TO INSURE A COMPLETELY ACCURATE DEFINITION OF THE BOUNDARY
C CONDITIONS AT THE CORNER MESH POINTS, THE USER SHOULD
C CONSULT THE APPROPRIATE DIFFERENCE APPROXIMATIONS (SEE
C SECTION 3 OF THE ACCOMPANYING PAPER).
C
C RETURN
C END
C
C 3) DIFFH AND DIFFV DEFINE FOR PDETWO THE HORIZONTAL AND VERTICAL
C DIFFUSION COEFFICIENTS RESPECTIVELY. THE CONSTRUCTION OF
C DIFFV IS ANALOGOUS TO THE CONSTRUCTION OF DIFFH. THE
C USER SHOULD CONSTRUCT DIFFH AS FOLLOWS ..
C
C SUBROUTINE DIFFH (T,X,Y,U,DH,NPDE)
C DIMENSION U(NPDE),DH(NPDE,NPDE)
C
C IN THIS ROUTINE DEFINE THE DH(L,K) COEFFICIENTS (L,K =
C 1,2,...,NPDE). THE INCOMING PARAMTERS X AND Y DENOTE EITHER
C A BOUNDARY POINT OR A MESH MIDPOINT.
C
C RETURN
C END
C
C 4) THE RIGHT HAND SIDE OF THE PDE IS DEFINED FOR PDETWO IN
C THE SUBROUTINE F. THIS ROUTINE IS CONSTRUCTED AS FOLLOWS..
C
C SUBROUTINE F (T,X,Y,U,UX,UY,DUXX,DUY,DUY,DUY,DUY,DUY,DUY,NPDE)
C DIMENSION U(NPDE), UX(NPDE), UY(NPDE), DUXX(NPDE,NPDE),
C * DUY(NPDE,NPDE), DUY(NPDE)
C
C IN THIS ROUTINE, THE INCOMING VALUES X AND Y REPRESENT
C THE MESH POINT BEING EVALUATED AND UX, UY, DUXX AND DUY
C ARE THE VALUES DENOTED IN THE PROBLEM DEFINITION ABOVE.
C USING THESE VALUES, DEFINE IN DUY(L) (L = 1,2,...,NPDE)
C THE RIGHT HAND SIDE OF THE PDE*S.
C
C RETURN
C END
C
C*****
C THE INPUT PARAMETERS
C*****
C NPDE IS THE NUMBER OF PARTIAL DIFFERENTIAL EQUATIONS.
C NX IS THE NUMBER OF MESH POINTS IN THE HORIZONTAL
C DIRECTION.
C NY IS THE NUMBER OF MESH POINTS IN THE VERTICAL
C DIRECTION.
C X ARE THE MESH POINTS IN THE HORIZONTAL DIRECTION.
C Y ARE THE MESH POINTS IN THE VERTICAL DIRECTION.
C T IS THE CURRENT TIME.
C U CONTAINS THE CURRENT SOLUTION VALUES FOR ALL THE
C MESH POINTS.

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C DVS          SAVES THE VERTICAL DIFFUSION COEFFICIENTS FOR
C              FUTURE EVALUATIONS.
C DVST        SAVES THE VERTICAL DIFFUSION COEFFICIENTS FOR THE
C              FUTURE EVALUATIONS OF THE TOP ROW OF MESH POINTS.
C DH          RETURNS FROM DIFFH THE HORIZONTAL DIFFUSION COEFFI-
C              CIENTS AND CONTAINS DUXX, THE APPROXIMATIONS TO THE
C              HORIZONTAL DIVERGENCE TERMS, ON CALLS TO F AND STORES
C              THE PREVIOUS VALUE OF DUXX.
C DV          RETURNS FROM DIFFV THE VERTICAL DIFFUSION COEFFI-
C              CIENTS AND CONTAINS DUYV, THE APPROXIMATIONS TO THE
C              VERTICAL DIVERGENCE TERMS, ON CALLS TO F.
C AH, BH, CH  CONTAIN THE HORIZONTAL BOUNDARY COEFFICIENTS
C              PASSED TO PDETWO FROM BNDRYH.
C AV, BV, CV  CONTAIN THE VERTICAL BOUNDARY COEFFICIENTS
C              PASSED TO PDETWO FROM BNDRYV.
C UX,UY       STORE THE DIFFERENCE APPROXIMATIONS FOR THE FIRST
C              PARTIAL OF U WITH RESPECT TO X AND Y RESPECTIVELY.
C DXI,DXIR,
C   DXIC      STORE INFORMATION ABOUT THE HORIZONTAL MESH SPACING.
C XAVG        STORES INFORMATION ABOUT THE AVERAGE BETWEEN TWO MESH
C              POINTS ON THE HORIZONTAL AXIS.
C UDVA,UDVB,
C   UDHR,UDHL STORE INFORMATION FOR APPROXIMATING DUXX AND DUYV.
C UAVGH,UAVGV CONTAIN U AVERAGES FOR APPROXIMATING THE DIFFUSION
C              COEFFICIENTS AT THE MESH MID-POINTS.
C
C*****
C THE OUTPUT PARAMETERS
C*****
C
C U           CONTAINS SOLUTIONS UPDATED TO TIME T FOR THE
C              BOUNDARY MESH POINTS DEFINED BY DIRICHLET BOUNDARY
C              CONDITIONS.
C DUDT       IS THE RIGHT HAND SIDE OF THE SYSTEM OF ODE*S PASSED
C              TO THE INTEGRATOR. THESE VALUES ARE CALCULATED
C              FROM THE CENTERED DIFFERENCE APPROXIMATIONS OF
C              THE SPATIAL VARIABLES.
C
C*****
C REFERENCES
C*****
C           1. A. C. HINDMARSH, GEARB.. SOLUTION OF ORDINARY DIFFERENTIAL
C              EQUATIONS HAVING BANDED JACOBIAN, UCID-30059 REV. 2,
C              LAWRENCE LIVERMORE LABORATORY, P.O.BOX 808, LIVERMORE,
C              CA 94550, JUNE 1977.
C
C           2. R.F. SINCOVEC AND N.K. MADSEN, SOFTWARE FOR NONLINEAR
C              PARTIAL DIFFERENTIAL EQUATIONS, ACM TRANSACTIONS ON
C              MATHEMATICAL SOFTWARE, SEPT. 1975, PP. 232-260.
C
C           3. R.F. SINCOVEC AND N.K. MADSEN, ALGORITHM 494 PDEONE,
C              SOLUTIONS OF SYSTEMS OF PARTIAL DIFFERENTIAL EQUATIONS,
C              ACM TRANSACTIONS ON MATHEMATICAL SOFTWARE, SEPT.
C              1975, PP. 262-263.
C
C*****
C           SUBROUTINE PSETM (NPDE,NX,NY,X,Y,U,UMAX,USAVE,DUDTR,DUDT,PW,CON,
C              *
C              MITER,IER,NEBAND,WORK,IWORK)
C*****
C
C           PSETM IS INTENDED TO BE USED IN CONJUNCTION WITH THE
C           INTERFACE PDETWO FOR THE SOLUTION OF SECOND ORDER TIME DEP-
C           ENDENT PARTIAL DIFFERENTIAL EQUATIONS (PDE*S) DEFINED OVER A
C           RECTANGULAR REGION. PSETM IS SPECIFICALLY DESIGNED TO REPLACE
C           THE ROUTINE PSETB USED IN THE ORDINARY DIFFERENTIAL EQUATION
C           INTEGRATOR GEARB (1) IN ORDER TO MINIMIZE THE COMPUTATIONS
C           REQUIRED TO GENERATE THE JACOBIAN MATRIX.
C

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C     PSETM IS CALLED BY STIFFP, A ROUTINE IN THE MODIFIED GEARB,
C     WHEN THE INTEGRATION METHOD (MITER=1 OR 2) REQUIRES A JACOBIAN
C     MATRIX. WHEN MITER=1, THE JACOBIAN IS DETERMINED BY A USER
C     DEFINED ROUTINE, PDB. THIS METHOD SHOULD BE AVOIDED SINCE IT RE-
C     QUIRES THE USER TO COMPLETELY UNDERSTAND PDETWO. IF THIS
C     METHOD IS USED, PSETM OFFERS NO ADVANTAGE OVER PSETB. IF
C     MITER=2, THE JACOBIAN IS APPROXIMATED BY FINITE DIFFERENCES,
C     USING THE APPROXIMATION  $J = (DUDT(U+R) - DUDT(U)) / R$ ,
C     WHERE DUDT(U+R) AND DUDT(U) ARE THE VALUES OF THE RIGHT HAND
C     SIDE OF THE PDE*S EVALUATED AT U+R AND U RESPECTIVELY, AND R
C     IS SOME SMALL NUMBER. PSETM TAKES ADVANTAGE OF THE STRUCTURE
C     OF THE JACOBIAN MATRIX REQUIRED BY PDETWO TO REDUCE THE COMPU-
C     TATIONS NEEDED TO GENERATE THE MATRIX (REQUIRING ONLY 5 * NPDE
C     CALLS TO PDETWO). WITH EACH CALL TO PDETWO, PSETM DETERMINES
C
C     THE ENTRIES IN THE JACOBIAN FOR THE MESH POINTS IN THE FOLLOWING
C     PATTERN ..
C
C           X O O O O X O O O O
C           O O O X O O O O X O
C           O X O O O O X O O O
C           O O O O X O O O O X
C           O O X O O O O X O O
C           X O O O O X O O O O
C
C     WHERE X REPRESENTS THE POINTS FOR WHICH THE JACOBIAN IS
C     APPROXIMATED. MITER=2 IS THE RECOMMENDED METHOD.
C
C     TO CREATE A DOUBLE PRECISION VERSION OF PSETM..
C     CHANGE THE REAL STATEMENT BELOW AND CHANGE THE SINGLE PRECISION
C     FUNCTIONS ABS, SQRT, AND AMAX1.
C
C*****
C     THE PARAMETERS
C*****
C
C     NPDE      IS THE NUMBER OF PARTIAL DIFFERENTIAL EQUATIONS.
C     NX        IS THE NUMBER OF MESH POINTS IN THE HORIZONTAL
C               DIRECTION.
C     NY        IS THE NUMBER OF MESH POINTS IN THE VERTICAL
C               DIRECTION.
C     X         ARE THE MESH POINTS IN THE HORIZONTAL DIRECTION.
C     Y         ARE THE MESH POINTS IN THE VERTICAL DIRECTION.
C     U         CONTAINS THE CURRENT SOLUTIONS OF THE PARTIAL
C               DIFFERENTIAL EQUATIONS.
C     UMAX      IS THE MAXIMUM U VALUES, WHICH ARE USED TO SCALE
C               THE VALUE IN R AND FOR ERROR CONTROL IN GEARB.
C     USAVE     IS A TEMPORARY STORE FOR U VALUES.
C     DUDTR     IS THE VALUE OF DUDT(U+R) RETURNED FROM PDETWO.
C     DUDT      IS THE VALUE OF DUDT(U), WHICH IS PASSED TO PSETM
C               FROM STIFFP.
C     PW        STORES THE APPROXIMATION FOR THE JACOBIAN.
C     CON       IS THE CONSTANT (-H*EL(1)).
C     MITER     INDICATES THE TYPE OF ITERATION METHOD BEING USED
C               BY THE INTEGRATOR.
C               MITER = 0 MEANS FUNCTIONAL ITERATION.
C               MITER = 1 MEANS THE CHORD METHOD WITH AN ANALYTIC
C                   JACOBIAN SUPPLIED IN THE USER DEFINED
C                   ROUTINE PDB. THIS METHOD IS IN GEARB,
C                   BUT IT SHOULD BE AVOIDED.
C               MITER = 2 MEANS THE CHORD METHOD WITH THE JACO-
C                   BIAN CALCULATED IN PSETM. THIS IS THE
C                   ONLY METHOD WHERE PSETM IS USEFUL.
C               MITER = 3 MEANS THE CHORD METHOD WITH THE JACO-
C                   BIAN REPLACED BY A DIAGONAL APPROXI-
C                   MATION BASED ON A DIRECTIONAL DERIV-
C                   ATIVE. THIS METHOD IS IN GEARB, BUT
C                   IS NOT GENERALLY USEFUL IN SOLVING PDE*S.
C     IER       IS AN ERROR INDICATOR USED IN THE ROUTINE DECBR.
C     NEBAND    = 3 * ML + 1

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C WORK          PROVIDES WORKING STORAGE FOR DYNAMIC DIMENSIONING
C              OF THE ARRAYS IN PDETWO AND THE MODIFIED GEARB.
C IWORK        IS THE PIVOT VECTOR USED IN THE LU DECOMPOSITION.
C
C*****
C THE VARIABLES
C*****
C R            IS A SMALL INCREMENT USED IN EVALUATING THE
C              FUNCTION NEAR THE CURRENT SOLUTION U ( I.E.
C              DUDT (U+R)).
C RØ          IS A LOWER BOUND ON THE SIZE OF R.
C D            IS A SMALL NUMBER USED TO APPROXIMATE THE
C              DERIVATIVE.
C NXNPDE      = NX * NPDE
C
C T            IS THE TIME BEING USED FOR THE INTEGRATION.
C H            IS THE CURRENT TIME STEP SIZE BEING USED IN THE
C              INTEGRATION.
C UROUND      IS THE UNIT ROUNDOFF OF THE MACHINE.
C EPSJ        IS SQRT(UROUND).
C ML          IS THE WIDTH OF THE LOWER (AND UPPER) HALF OF THE
C              BAND OF THE JACOBIAN.  ML = (NX + 1) * NPDE - 1.
C IW1 - IW27  ARE THE SUBSCRIPT POINTERS USED TO INDICATE THE
C              STORAGE LOCATIONS IN WORK OF ARRAYS IN PDETWO AND
C              THE MODIFIED GEARB.
C
C*****
C THE ROUTINES CALLED BY PSETM
C*****
C PDB(NODE,T,U,PW,NODE,ML,ML) IS A USER DEFINED ROUTINE WHICH DEFINES
C              THE JACOBIAN IN PW IF MITER=1.  SINCE THIS METHOD IS
C              NOT RECOMMENDED, A DUMMY ROUTINE IS PROVIDED.
C DECBR(NEBAND,NODE,ML,ML,PW,IWORK,IER) COMPUTES THE LU DECOMPOSITION
C              OF THE JACOBIAN FOR THE DIRECT SOLUTION OF THE
C              JACOBIAN.
C PDETWO(NPDE,NX,NY,X,Y,T,U,...) IS THE INTERFACE WHICH DESCRETIZES
C              THE SPATIAL VARIABLES OF THE TWO DIMENSIONAL SYSTEM OF
C              PDE*S, THEREBY EVALUATING THE RIGHT HAND SIDE OF
C              THE SYSTEM OF ODE*S.
C
C*****
C REFERENCES
C*****
C          1.  A. C. HINDMARSH, GEARB.. SOLUTION OF ORDINARY DIFFERENTIAL
C              EQUATIONS HAVING BANDED JACOBIAN, UCID-3ØØ59 REV. 2,
C              LAWRENCE LIVERMORE LABORATORY, P.O.BOX 8Ø8, LIVERMORE,
C              CA 9455Ø, JUNE 1977.
C
C*****

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ALGORITHM 566

FORTRAN Subroutines for Testing Unconstrained Optimization Software [C5], [E4]

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Key Words and Phrases: performance testing, systems of nonlinear equations, nonlinear least squares, unconstrained minimization, optimization software
CR Categories: 4.6, 5.15, 5.41
Language: FORTRAN

DESCRIPTION

This is the FORTRAN package of subroutines described in [1] for testing unconstrained optimization software. The following three problem areas are considered.

- (1) Zeros of systems of N nonlinear functions in N variables.
- (2) Least squares minimization of M nonlinear functions in N variables.
- (3) Unconstrained minimization of an objective function with N variables.

The subroutines that define the test functions and starting points depend on the dimension parameters M and N and on the problem number $NPROB$. We first describe the subroutines for the test functions.

For systems of nonlinear functions,

VECFCN(N , X , $FVEC$, $NPROB$)

returns the function values in the N -vector $FVEC$, and

VECJAC(N , X , $FJAC$, $LDFJAC$, $NPROB$)

returns the Jacobian matrix in the N by N array $FJAC$. (The parameter $LDFJAC$ is the leading dimension of the array $FJAC$ as defined in the main program.) In order to prevent gross inefficiencies with solvers that only require one function value at a time,

COMFCN(N , K , X , $FCNK$, $NPROB$)

returns the K th function value in $FCNK$.

For nonlinear least squares,

SSQFCN(M , N , X , $FVEC$, $NPROB$)

returns the function values in the M -vector $FVEC$, and

SSQJAC(M , N , X , $FJAC$, $LDFJAC$, $NPROB$)

returns the Jacobian matrix in the M by N array $FJAC$.

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For unconstrained minimization,

OBJFCN(N, X, F, NPROB)

returns the objective function value in F, and

GRDFCN(N, X, G, NPROB)

returns the gradient components in the N-vector G.

For each problem area, the starting points are generated by

INITPT(N, X, NPROB, FACTOR)

which returns in X the starting point corresponding to the parameters NPROB and FACTOR. If XS denotes the standard starting point, then X will contain FACTOR*XS, except that if XS is the zero vector and FACTOR is not unity, then all the components of X will be set to FACTOR.

To test a code in any of the three problem areas, the user must provide a driver and interface routine. The driver reads in the data that define the dimensions, the problem number, and FACTOR, calls INITPT, and then calls the code of interest and prints out results. The interface routine provides a link between the code with its particular function routine calling sequences and the subroutines for the test functions.

The package includes example drivers and interface routines for each of the problem areas. Sample data are also provided.

REFERENCES

1. MORÉ, J.J., GARBOW, B.S., AND HILLSTROM, K.E. Testing unconstrained optimization software. *ACM Trans. Math. Soft.* 7, 1 (March 1981), 17-41.

ALGORITHM

[Part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service].

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C      THIS PROGRAM TESTS CODES FOR THE UNCONSTRAINED OPTIMIZATION OF
C      A NONLINEAR FUNCTION OF N VARIABLES. IT CONSISTS OF A DRIVER
C      AND AN INTERFACE SUBROUTINE FCN. THE DRIVER READS IN DATA,
C      CALLS THE UNCONSTRAINED OPTIMIZER, AND FINALLY PRINTS OUT
C      INFORMATION ON THE PERFORMANCE OF THE OPTIMIZER. THIS IS
C      ONLY A SAMPLE DRIVER, MANY OTHER DRIVERS ARE POSSIBLE. THE
C      INTERFACE SUBROUTINE FCN IS NECESSARY TO TAKE INTO ACCOUNT THE
C      FORMS OF CALLING SEQUENCES USED BY THE FUNCTION SUBROUTINES
C      IN THE VARIOUS UNCONSTRAINED OPTIMIZERS.
C
C      SUBPROGRAMS CALLED
C
C      USER-SUPPLIED ..... ENORM,FCN,SOLVER
C
C      MINPACK-SUPPLIED ... GRDFCN,INITPT,OBJFCN
C
C      FORTRAN-SUPPLIED ... DSQRT
C
C      MINPACK. VERSION OF NOVEMBER 1978.
C      BURTON S. GARBOW, KENNETH E. HILLSTROM, JORGE J. MORE
C      SUBROUTINE FCN(N,X,F,GVEC,IFLAG)
C      INTEGER N,IFLAG
C      DOUBLE PRECISION F
C      DOUBLE PRECISION X(N),GVEC(N)
C      *****
C
C      THE CALLING SEQUENCE OF FCN SHOULD BE IDENTICAL TO THE
C      CALLING SEQUENCE OF THE FUNCTION SUBROUTINE IN THE
C      UNCONSTRAINED OPTIMIZER. FCN SHOULD ONLY CALL THE TESTING
C      FUNCTION AND GRADIENT SUBROUTINES OBJFCN AND GRDFCN WITH
C      THE APPROPRIATE VALUE OF PROBLEM NUMBER (NPROB).
C

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C   SUBPROGRAMS CALLED
C
C   MINPACK-SUPPLIED ... GRDFCN,OBJFCN
C
C   MINPACK. VERSION OF JULY 1978.
C   BURTON S. GARBOW, KENNETH E. HILLSTROM, JORGE J. MORE
C
C   *****
C   INTEGER NPROB,NFEV
C   COMMON /REFNUM/ NPROB,NFEV
C   CALL OBJFCN(N,X,F,NPROB)
C   CALL GRDFCN(N,X,GVEC,NPROB)
C   NFEV = NFEV + 1
C   RETURN
C
C   LAST CARD OF INTERFACE SUBROUTINE FCN.
C
C   END
C   SUBROUTINE INITPT(N,X,NPROB,FACTOR)
C   INTEGER N,NPROB
C   DOUBLE PRECISION FACTOR
C   DOUBLE PRECISION X(N)
C   *****
C
C   SUBROUTINE INITPT
C
C   THIS SUBROUTINE SPECIFIES THE STANDARD STARTING POINTS FOR THE
C   FUNCTIONS DEFINED BY SUBROUTINE OBJFCN. THE SUBROUTINE RETURNS
C   IN X A MULTIPLE (FACTOR) OF THE STANDARD STARTING POINT. FOR
C   THE SEVENTH FUNCTION THE STANDARD STARTING POINT IS ZERO, SO IN
C   THIS CASE, IF FACTOR IS NOT UNITY, THEN THE SUBROUTINE RETURNS
C   THE VECTOR X(J) = FACTOR, J=1,...,N.
C
C   THE SUBROUTINE STATEMENT IS
C
C   SUBROUTINE INITPT(N,X,NPROB,FACTOR)
C
C   WHERE
C
C   N IS A POSITIVE INTEGER INPUT VARIABLE.
C
C   X IS AN OUTPUT ARRAY OF LENGTH N WHICH CONTAINS THE STANDARD
C   STARTING POINT FOR PROBLEM NPROB MULTIPLIED BY FACTOR.
C
C   NPROB IS A POSITIVE INTEGER INPUT VARIABLE WHICH DEFINES THE
C   NUMBER OF THE PROBLEM. NPROB MUST NOT EXCEED 18.
C
C   FACTOR IS AN INPUT VARIABLE WHICH SPECIFIES THE MULTIPLE OF
C   THE STANDARD STARTING POINT. IF FACTOR IS UNITY, NO
C   MULTIPLICATION IS PERFORMED.
C
C   MINPACK. VERSION OF JULY 1978.
C   BURTON S. GARBOW, KENNETH E. HILLSTROM, JORGE J. MORE
C   SUBROUTINE OBJFCN(N,X,F,NPROB)
C   INTEGER N,NPROB
C   DOUBLE PRECISION F
C   DOUBLE PRECISION X(N)
C   *****
C
C   SUBROUTINE OBJFCN
C
C   THIS SUBROUTINE DEFINES THE OBJECTIVE FUNCTIONS OF EIGHTEEN
C   NONLINEAR UNCONSTRAINED MINIMIZATION PROBLEMS. THE VALUES
C   OF N FOR FUNCTIONS 1,2,3,4,5,10,11,12,16 AND 17 ARE
C   3,6,3,2,3,2,4,3,2 AND 4, RESPECTIVELY.
C   FOR FUNCTION 7, N MAY BE 2 OR GREATER BUT IS USUALLY 6 OR 9.
C   FOR FUNCTIONS 6,8,9,13,14,15 AND 18 N MAY BE VARIABLE,
C   HOWEVER IT MUST BE EVEN FOR FUNCTION 14, A MULTIPLE OF 4 FOR
C   FUNCTION 15, AND NOT GREATER THAN 50 FOR FUNCTION 18.
C

```

```
C THE SUBROUTINE STATEMENT IS
C
C SUBROUTINE OBJFCN(N,X,F,NPROB)
C
C WHERE
C
C N IS A POSITIVE INTEGER INPUT VARIABLE.
C
C X IS AN INPUT ARRAY OF LENGTH N.
C
C F IS AN OUTPUT VARIABLE WHICH CONTAINS THE VALUE OF
C THE NPROB OBJECTIVE FUNCTION EVALUATED AT X.
C
C NPROB IS A POSITIVE INTEGER INPUT VARIABLE WHICH DEFINES THE
C NUMBER OF THE PROBLEM. NPROB MUST NOT EXCEED 18.
C
C SUBPROGRAMS CALLED
C
C FORTRAN-SUPPLIED ... DABS,DATAN,DCOS,DEXP,DLOG,DSIGN,DSIN,
C DSQRT
C
C MINPACK. VERSION OF JULY 1978.
C BURTON S. GARBOW, KENNETH E. HILLSTROM, JORGE J. MORE
C SUBROUTINE GRDFCN(N,X,G,NPROB)
C INTEGER N,NPROB
C DOUBLE PRECISION X(N),G(N)
C *****
C
C SUBROUTINE GRDFCN
C
C THIS SUBROUTINE DEFINES THE GRADIENT VECTORS OF EIGHTEEN
C NONLINEAR UNCONSTRAINED MINIMIZATION PROBLEMS. THE PROBLEM
C DIMENSIONS ARE AS DESCRIBED IN THE PROLOGUE COMMENTS OF OBJFCN.
C
C THE SUBROUTINE STATEMENT IS
C
C SUBROUTINE GRDFCN(N,X,G,NPROB)
C
C WHERE
C
C N IS A POSITIVE INTEGER INPUT VARIABLE.
C
C X IS AN INPUT ARRAY OF LENGTH N.
C
C G IS AN OUTPUT ARRAY OF LENGTH N WHICH CONTAINS THE COMPONENTS
C OF THE GRADIENT VECTOR OF THE NPROB OBJECTIVE FUNCTION
C EVALUATED AT X.
C
C NPROB IS A POSITIVE INTEGER INPUT VARIABLE WHICH DEFINES THE
C NUMBER OF THE PROBLEM. NPROB MUST NOT EXCEED 18.
C
C SUBPROGRAMS CALLED
C
C FORTRAN-SUPPLIED ... DABS,DATAN,DCOS,DEXP,DLOG,DSIGN,DSIN,
C DSQRT
C
C MINPACK. VERSION OF JULY 1978.
C BURTON S. GARBOW, KENNETH E. HILLSTROM, JORGE J. MORE
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ALGORITHM 567

Extended-Range Arithmetic and Normalized Legendre Polynomials [A1], [C1]

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and

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Key Words and Phrases: angular momentum, extended-range arithmetic, Legendre polynomials, overflow, underflow

CR Categories: 3.17, 4.9, 5.12

Language: FORTRAN

DESCRIPTION

This algorithm consists of two logically distinct parts: (1) a package of six FORTRAN subroutines to facilitate the use of a special form of computer floating-point arithmetic that we call extended-range arithmetic; and (2) a FORTRAN subroutine that computes values of normalized Legendre polynomials according to an algorithm that generates (for some inputs) floating-point numbers that are outside the range of any computer. Our desire to produce a robust FORTRAN subroutine to compute these polynomials stimulated the development of the extended-range software package. This package may prove to be useful for many other computations.

Normalized Legendre polynomials are defined by the formula

$$\bar{P}_\nu^\mu(x) = \left\{ \left(\nu + \frac{1}{2} \right) \frac{(\nu - \mu)!}{(\nu + \mu)!} \right\}^{1/2} (1 - x^2)^{\mu/2} \frac{d^\mu}{dx^\mu} P_\nu(x),$$

where μ and ν are nonnegative integers, x is a real variable lying in the closed interval $[-1, 1]$, and $P_\nu(x)$ is the ordinary Legendre polynomial. These functions satisfy a three-term recurrence relation in μ that is useful in computing $\bar{P}_\nu^\mu(x)$ for fixed ν and x , and sequences $\mu_1, \mu_1 + 1, \dots, \mu_2$ of values of μ . For stability reasons, the recurrence is applied in the backward direction, starting at $\mu = \nu + 1$ and $\mu = \nu$ and proceeding through $\nu - 1, \nu - 2, \dots, \mu_1$. The starting value of $\bar{P}_\nu^\mu(x)$ is determined from a first-order recurrence relation, and $\bar{P}_\nu^{\nu+1}(x) = 0$ for all ν, x . When x is close to ± 1 and ν is moderately large, this method fails because of underflow of early values in the recurrence sequence. Nevertheless, the method is attractive because of its inherent stability and simplicity.

A simple extension of floating-point arithmetic, called extended-range arithmetic, overcomes this difficulty. A real number ξ is represented in three storage locations. Two of these hold a number x in ordinary double-precision form; the third holds a signed integer k . The value of ξ is given by

$$\xi = x \times r^k,$$

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where r is the base (or radix) of double-precision numbers. In this (nonunique) representation x is called the *principal part* and k is called the *auxiliary index*. Addition, subtraction, multiplication, and division of extended-range numbers are, of course, trivial to define. The main challenge was to design subroutines that are portable and, at the same time, as efficient as possible. Addition and conversion of extended-range numbers to a decimal form suitable for printing proved to be the most complicated. The subroutine for addition uses nine cases to preclude underflow and overflow of the principal part of the result. The algorithm used in conversion to decimal form requires triple-precision calculation to achieve double-precision accuracy in the converted value of r^k , because of the possible large size of the auxiliary index k .

ACKNOWLEDGMENT

The authors are pleased to acknowledge the aid of R. N. Freemire of the National Bureau of Standards in testing the software and critically reviewing the FORTRAN coding.

REFERENCES

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ALGORITHM

[Usage of the extended-range software is described in the comments in the initializing subroutine SETUP. These are reproduced below. Similarly, usage of NORMP (the subroutine for computing normalized Legendre polynomials) is described below. The complete listings are available from the ACM Algorithms Distribution Service].

```

C *** SUBROUTINE NORMP ***
C   SUBROUTINE NORMP(NU, MU1, MU2, ARG, MODE, PN, IPN, ISIG)
C   INTEGER NU, MU1, MU2, MODE, IPN, ISIG
C   DOUBLE PRECISION ARG, PN
C
C   SUBROUTINE TO CALCULATE NORMALIZED LEGENDRE POLYNOMIALS
C   OF VARYING ORDER AND OF FIXED ARGUMENT AND DEGREE.
C   THE ORDER MU AND DEGREE NU ARE NONNEGATIVE INTEGERS AND THE
C   ARGUMENT IS REAL. THE ALGORITHM REQUIRES THE USE OF
C   NUMBERS OUTSIDE THE NORMAL MACHINE RANGE. THEREFORE,
C   THIS SUBROUTINE EMPLOYS A SPECIAL ARITHMETIC CALLED
C   EXTENDED-RANGE ARITHMETIC. (SEE SMITH, J.M., OLVER,
C   F.W.J., AND LOZIER, D.W., EXTENDED-RANGE ARITHMETIC
C   AND NORMALIZED LEGENDRE POLYNOMIALS, ACM TRANSACTIONS
C   ON MATHEMATICAL SOFTWARE, 1981).
C
C   IN EXTENDED-RANGE ARITHMETIC EACH NUMBER IS REPRESENTED
C   AS A PAIR (X,IX) WHICH HAS THE VALUE X*RADIX**IX, WHERE
C   RADIX IS THE BASE IN WHICH CALCULATIONS ARE PERFORMED. FOR
C   A FULL DESCRIPTION OF EXTENDED-RANGE ARITHMETIC, SEE THE
C   COMMENTS AT THE BEGINNING OF SUBROUTINE SETUP.
C
C   THE NORMALIZED LEGENDRE POLYNOMIAL IS A MULTIPLE OF THE
C   ASSOCIATED LEGENDRE POLYNOMIAL OF THE FIRST KIND
C   WHERE THE NORMALIZING COEFFICIENT IS CHOSEN SO THAT
C   THE INTEGRAL OF THE SQUARE OF THE FUNCTION FROM -1
C   TO +1 IS EQUAL TO 1 (SEE JAHNKE, E. EMDE, F. AND LOSCH, F.,
C   TABLES OF HIGHER FUNCTIONS, MCGRAW HILL, NEW YORK,
C   1960, P. 121).
C
C   THE INPUT VALUES TO NORMP ARE NU, MU1, MU2, ARG, AND MODE.
C   THESE MUST SATISFY:
C   1. NU IS A NON-NEGATIVE INTEGER SPECIFYING THE DEGREE.
C   2. MU1 AND MU2 ARE NON-NEGATIVE INTEGERS, WITH
C      MU1.LE.MU2, SPECIFYING THE RANGE OF ORDERS.
C   3. MODE IS AN INTEGER, RESTRICTED TO 1 OR 2,
C      WHICH INDICATES WHICH OF TWO FORMS OF THE

```

```

C          DOUBLE-PRECISION VARIABLE ARG IS INTENDED:
C          A. IF MODE=1 THEN NORMALIZED LEGENDRE(NU,MU,ARG) IS
C             COMPUTED FOR ALL MU SUCH THAT MU1 .LE. MU .LE. MU2 .
C             IN THIS CASE, -1 .LE. ARG .LE. 1 MUST BE SATISFIED.
C          B. IF MODE=2 THEN NORMALIZED LEGENDRE(NU,MU,COS ARG) IS
C             COMPUTED FOR ALL MU SUCH THAT MU1 .LE. MU .LE. MU2 ,
C             IN THIS CASE, -PI .LT. ARG .LT. PI MUST BE SATISFIED.
C
C          THE OUTPUT OF SUBROUTINE NORMP CONSISTS OF THE TWO
C          ARRAYS PN AND IPN AND THE ERROR ESTIMATE ISIG. THE
C          NORMALIZED LEGENDRE POLYNOMIALS ARE STORED AS EXTENDED-
C          RANGE NUMBERS WITH DOUBLE-PRECISION PRINCIPAL PARTS
C          IN ARRAY PN AND AUXILIARY INDICES IN ARRAY IPN.  THUS
C          (PN(1),IPN(1))=NORMALIZED LEGENDRE(NU,MU1,X)
C          (PN(2),IPN(2))=NORMALIZED LEGENDRE(NU,MU1+1,X)
C          .
C          .
C          (PN(K),IPN(K))=NORMALIZED LEGENDRE(NU,MU2,X)
C          WHERE K=MU2-MU1+1 AND X=ARG IF MODE=1, OR X=COS ARG IF MODE=2.
C          FINALLY, ISIG IS AN ESTIMATE OF THE NUMBER OF DECIMAL DIGITS
C          LOST THROUGH ROUNDING ERRORS IN THE COMPUTATION.  IF ARG IS
C          ACCURATE TO N SIGNIFICANT DECIMALS, THEN THE COMPUTED FUNCTION
C          VALUES ARE ACCURATE TO N-ISIG SIGNIFICANT DECIMALS (EXCEPT IN
C          NEIGHBORHOODS OF ZEROS OF THE FUNCTIONS).
C *** SUBROUTINE SETUP ***
C USAGE
C          CALL SETUP(IRAD,NRADPL,DZERO,NBITS)
C DESCRIPTION
C          SUBROUTINE SETUP MUST BE CALLED PRIOR TO CALLING ANY OTHER
C          EXTENDED-RANGE SUBROUTINE.  IT CALCULATES AND STORES SEVERAL
C          MACHINE-DEPENDENT CONSTANTS IN COMMON BLOCKS.  THE USER MUST
C          SUPPLY FOUR CONSTANTS THAT PERTAIN TO HIS PARTICULAR COMPUTER.
C          THE CONSTANTS ARE
C
C          IRAD = THE INTERNAL BASE OF DOUBLE-PRECISION
C                 ARITHMETIC IN THE COMPUTER.
C          NRADPL = THE NUMBER OF RADIX PLACES CARRIED IN
C                 THE DOUBLE-PRECISION REPRESENTATION.
C          DZERO = THE SMALLEST OF 1/DMIN, DMAX, DMAXLN WHERE
C                 DMIN = THE SMALLEST POSITIVE DOUBLE-PRECISION
C                 NUMBER OR AN UPPER BOUND TO THIS NUMBER,
C                 DMAX = THE LARGEST DOUBLE-PRECISION NUMBER
C                 OR A LOWER BOUND TO THIS NUMBER,
C                 DMAXLN = THE LARGEST DOUBLE-PRECISION NUMBER
C                 SUCH THAT DLOG(DMAXLN) CAN BE COMPUTED BY THE
C                 FORTRAN SYSTEM (ON MOST SYSTEMS DMAXLN = DMAX).
C          NBITS = THE NUMBER OF BITS (EXCLUSIVE OF SIGN) IN
C                 AN INTEGER WORD.
C
C          SUBROUTINE SETUP(IRAD, NRADPL, DZERO, NBITS)
C          INTEGER IRAD, NRADPL, NBITS
C          DOUBLE PRECISION DZERO
C
C          THIS IS THE SETTING-UP SUBROUTINE FOR A PACKAGE OF SUBROUTINES
C          THAT FACILITATE THE USE OF EXTENDED-RANGE ARITHMETIC.  EXTENDED-RANGE
C          ARITHMETIC ON A PARTICULAR COMPUTER IS DEFINED ON THE SET OF NUMBERS
C          OF THE FORM
C
C          (X,IX) = X*RADIX**IX
C
C          WHERE X IS A DOUBLE-PRECISION NUMBER CALLED THE PRINCIPAL PART,
C          IX IS AN INTEGER CALLED THE AUXILIARY INDEX, AND RADIX IS THE
C          INTERNAL BASE OF THE DOUBLE-PRECISION ARITHMETIC.  OBVIOUSLY,
C          EACH REAL NUMBER IS REPRESENTABLE WITHOUT ERROR BY MORE THAN ONE
C          EXTENDED-RANGE FORM.  CONVERSIONS BETWEEN DIFFERENT FORMS ARE
C          ESSENTIAL IN CARRYING OUT ARITHMETIC OPERATIONS.  WITH THE CHOICE
C          OF RADIX WE HAVE MADE, AND THE SUBROUTINES WE HAVE WRITTEN, THESE
C          CONVERSIONS ARE PERFORMED WITHOUT ERROR (AT LEAST ON MOST COMPUTERS).
C          (SEE SMITH, J.M., OLVER, F.W.J., AND LOZIER, D.W., EXTENDED-RANGE
C          ARITHMETIC AND NORMALIZED LEGENDRE POLYNOMIALS, ACM TRANSACTIONS ON
C          MATHEMATICAL SOFTWARE, 1981).
C
C          AN EXTENDED-RANGE NUMBER (X,IX) IS SAID TO BE IN ADJUSTED FORM IF
C          X AND IX ARE ZERO OR
C
C          RADIX**(-L) .LE. DABS(X) .LT. RADIX**L
C

```

```

C IS SATISFIED, WHERE L IS A COMPUTER-DEPENDENT INTEGER DEFINED IN THIS
C SUBROUTINE. TWO EXTENDED-RANGE NUMBERS IN ADJUSTED FORM CAN BE ADDED,
C SUBTRACTED, MULTIPLIED OR DIVIDED (IF THE DIVISOR IS NONZERO) WITHOUT
C CAUSING OVERFLOW OR UNDERFLOW IN THE PRINCIPAL PART OF THE RESULT.
C WITH PROPER USE OF THE EXTENDED-RANGE SUBROUTINES, THE ONLY OVERFLOW
C THAT CAN OCCUR IS INTEGER OVERFLOW IN THE AUXILIARY INDEX. IF THIS
C IS DETECTED, THE EXTENDED-RANGE SOFTWARE PRINTS A MESSAGE AND STOPS.
C
C   MULTIPLICATION AND DIVISION IS PERFORMED BY SETTING
C
C           (X,IX)*(Y,IY) = (X*Y,IX+IY)
C OR
C           (X,IX)/(Y,IY) = (X/Y,IX-IY).
C
C PRE-ADJUSTMENT OF THE OPERANDS IS ESSENTIAL TO AVOID
C OVERFLOW OR UNDERFLOW OF THE PRINCIPAL PART. SUBROUTINE
C ADJUST (SEE BELOW) MAY BE CALLED TO TRANSFORM ANY EXTENDED-
C RANGE NUMBER INTO ADJUSTED FORM.
C
C   ADDITION AND SUBTRACTION REQUIRE THE USE OF SUBROUTINE ADD
C (SEE BELOW). THE INPUT OPERANDS NEED NOT BE IN ADJUSTED FORM.
C HOWEVER, THE RESULT OF ADDITION OR SUBTRACTION IS RETURNED
C IN ADJUSTED FORM. THUS, FOR EXAMPLE, IF (X,IX),(Y,IY),
C (U,IU), AND (V,IV) ARE IN ADJUSTED FORM, THEN
C
C           (X,IX)*(Y,IY) + (U,IU)*(V,IV)
C
C CAN BE COMPUTED AND STORED IN ADJUSTED FORM WITH NO EXPLICIT
C CALLS TO ADJUST.
C
C   WHEN AN EXTENDED-RANGE NUMBER IS TO BE PRINTED, IT MUST BE
C CONVERTED TO AN EXTENDED-RANGE FORM WITH DECIMAL RADIX. SUBROUTINE
C CNVRT IS PROVIDED FOR THIS PURPOSE.
C
C   THE SUBROUTINES CONTAINED IN THIS PACKAGE ARE
C
C *** SUBROUTINE ADD ***
C USAGE
C           CALL ADD(X,IX,Y,IY,Z,IZ)
C DESCRIPTION
C           FORMS THE EXTENDED-RANGE SUM (Z,IZ) =
C           (X,IX) + (Y,IY). (Z,IZ) IS ADJUSTED
C           BEFORE RETURNING. THE INPUT OPERANDS
C           NEED NOT BE IN ADJUSTED FORM, BUT THEIR
C           PRINCIPAL PARTS MUST SATISFY
C           RADIX**(-2L).LE.DABS(X).LE.RADIX**(2L),
C           RADIX**(-2L).LE.DABS(Y).LE.RADIX**(2L).
C
C *** SUBROUTINE ADJUST ***
C USAGE
C           CALL ADJUST(X,IX)
C DESCRIPTION
C           TRANSFORMS (X,IX) SO THAT
C           RADIX**(-L).LE.DABS(X).LT.RADIX**L.
C           ON MOST COMPUTERS THIS TRANSFORMATION DOES
C           NOT CHANGE THE MANTISSA OF X PROVIDED RADIX IS
C           THE NUMBER BASE OF DOUBLE-PRECISION ARITHMETIC.
C
C *** SUBROUTINE CNV210 ***
C USAGE
C           CALL CNV210(K,Z,J)
C DESCRIPTION
C           GIVEN K THIS SUBROUTINE COMPUTES J AND Z
C           SUCH THAT RADIX**K = Z*10**J, WHERE Z IS IN
C           THE RANGE 1/10 .LE. Z .LT. 1.
C           THE VALUE OF Z WILL BE ACCURATE TO FULL
C           DOUBLE PRECISION PROVIDED THE NUMBER
C           OF DECIMAL PLACES IN THE LARGEST
C           INTEGER PLUS THE NUMBER OF DECIMAL
C           PLACES CARRIED IN DOUBLE PRECISION DOES NOT
C           EXCEED 60. CNV210 IS CALLED BY SUBROUTINE
C           CNVRT WHEN NECESSARY. THE USER SHOULD
C           NEVER NEED TO CALL CNV210 DIRECTLY.
C
C *** SUBROUTINE CNVRT ***
C USAGE
C           CALL CNVRT(X,IX)

```

```

C DESCRIPTION
C          CONVERTS (X,IX) = X*RADIX**IX
C          TO DECIMAL FORM IN PREPARATION FOR
C          PRINTING, SO THAT (X,IX) = X*10**IX
C          WHERE 1/10 .LE. DABS(X) .LT. 1
C          IS RETURNED, EXCEPT THAT IF
C          (DABS(X),IX) IS BETWEEN RADIX**(-2L)
C          AND RADIX**(2L) THEN THE REDUCED
C          FORM WITH IX = 0 IS RETURNED.
C
C *** SUBROUTINE REDUCE ***
C USAGE
C          CALL REDUCE(X,IX)
C DESCRIPTION
C          IF
C          RADIX**(-2L) .LE. (DABS(X),IX) .LE. RADIX**(2L)
C          THEN REDUCE TRANSFORMS (X,IX) SO THAT IX=0.
C          IF (X,IX) IS OUTSIDE THE ABOVE RANGE,
C          THEN REDUCE TAKES NO ACTION.
C          THIS SUBROUTINE IS USEFUL IF THE
C          RESULTS OF EXTENDED-RANGE CALCULATIONS
C          ARE TO BE USED IN SUBSEQUENT ORDINARY
C          DOUBLE-PRECISION CALCULATIONS.
C
C          LUERR IS THE OUTPUT UNIT NUMBER FOR ERROR MESSAGES, SET
C          EQUAL TO 6 BELOW. SETUP DETERMINES THE OTHER COMMON VARIABLES
C          FROM THE INPUT.
C          COMMON /EXR1/ LUERR
C          DOUBLE PRECISION RADIX, RADIXL, RAD2L, DLG10R
C          COMMON /EXR2/ RADIX, RADIXL, RAD2L, DLG10R, L, L2, KMAX
C          COMMON /EXR3/ NLG102, MLG102, LG102(21)

```


ALGORITHM 568

PDS—A Portable Directory System

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Key Words and Phrases: file system, UNIX, RATFOR
CR Categories: 4.19, 4.33, 4.35, 4.41
Language: RATFOR (FORTRAN)

1. DESCRIPTION

PDS is a set of procedures that provides a machine-independent method of file specification. PDS provides capabilities beyond those provided by many vendor-supplied systems. In addition, because PDS is portable, additional capabilities, such as protection schemes, file usage statistics, or some of the functions of a source code control system [5], can be added easily.

The basic function of PDS is to maintain a useful directory structure and provide a set of primitives for manipulating that structure. The PDS directory structure is identical to the tree structure of the UNIX [4] file system, and many of the PDS primitives are identical to UNIX primitives. PDS is, in large part, a portable implementation of the UNIX directory system. Besides PDS's machine-independence, the major differences are the extensibility of PDS and, as described in the next section, its i/o independence.

In the simplest terms, PDS provides a directory structure and a *mapping* from machine-independent file names to machine-dependent names. It deals only with the information *describing* a file; it does not use or manipulate actual files in any way. The importance of this approach is that PDS is used to specify a file but does not participate in the actual i/o to that file. Consequently, there is no impact on i/o efficiency when PDS is used.

PDS manipulates a rooted tree structure in which the leaves are files or directories and the nodes are directories. A directory is simply a list of files and directories. An example is shown in Figure 1, in which circles indicate directories and squares indicate files. The root of the tree is denoted by "/", and files and directories are denoted by their "path," which specifies their absolute position in the tree. A path is composed of the names of the nodes on the path from the root to the desired file or directory. The path components are separated by slashes; for example, the path for file "file2" in Figure 1 is "/m/n/file2." The directory entries "." and ".." refer, respectively, to the directory itself and to the immediate ancestor. These names may be used as path components, providing an explicit means of using the structural properties of the tree. If a path does not begin with "/", it is taken to be rooted at the "current directory." For example, with the

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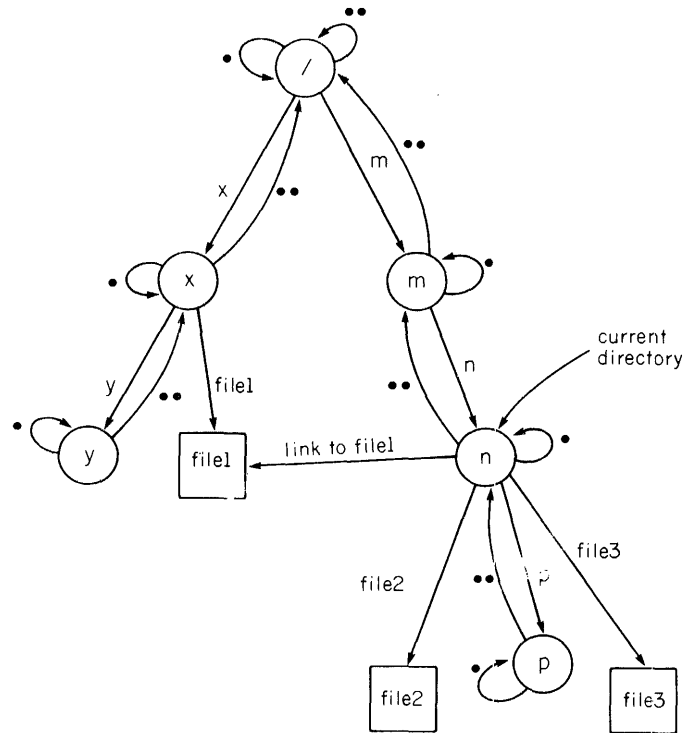


Fig. 1. A directory system.

current directory at “n” in Figure 1, the path “file3” is equivalent to “/m/n/file3.”

Files and directories are equivalent with the exception that directories are manipulated by PDS and files are not. Files simply “contain” the machine-dependent names of the actual files. These machine-dependent names are referred to as *host names*. The basic function of PDS is to map paths into host names.

There are ten PDS primitives; they are summarized in Table I. Detailed descriptions of each primitive are given in [1] and in the machine-readable comments. The most important primitives are *openf* and *creatf*, which map a path to a host name and perform machine-dependent operations concerned with preparing files for i/o. The rest are concerned primarily with manipulating the directory structure.

openf obtains the host name corresponding to *path* and calls a machine-dependent routine to actually open the file. *mode* indicates how the file is to be opened (e.g., “read”, “write”, “append”, etc.). *mode* is not modified by *openf*, so it can be whatever is most appropriate on the host system. *openf* returns what the machine-dependent open routine returns, which is whatever is most useful as an argument in calling machine-dependent i/o routines that perform the actual i/o (e.g., **read** and **write**). Typically, a channel number or FORTRAN unit number is returned.

creatf is similar to *openf*, except that *path* is created along with a host name and file. The generation of the host name is performed automatically. After the file is created, it is opened according to *mode* as if *openf* had been called.

openf and *creatf* are insensitive to the values of *mode*. It is, however, useful to introduce standard values representing common file usage, such as reading and writing. This approach promotes machine-independence in programs that use PDS and is used, for example, with the programs described in [3]. Such *mode* values are scrutinized by the machine-dependent i/o interface routines, however, permitting PDS to be used in almost any environment—not just those in which i/o capabilities conform to preconceived notions of “common” file usage.

PDS has *no* other i/o primitives; it is completely independent of the actual i/o. Thus, additional overhead is incurred in opening files, but none is incurred in accessing them. The effect of the additional overhead is minimal since the former

Table I. PDS Primitives

<i>chdir(path)</i>	change current directory to <i>path</i>
<i>chds(root, ilst)</i>	change to another directory system
<i>creatf(path, mode)</i>	create <i>path</i> and open it for i/o
<i>link(path1, path2)</i>	make a link to <i>path1</i> named <i>path2</i>
<i>mkdir(path)</i>	make a directory named <i>path</i>
<i>mkds(root, ilst)</i>	make a directory system
<i>mkfile(path, hname)</i>	make a file <i>path</i> with host name <i>hname</i>
<i>openf(path, mode)</i>	open <i>path</i> for i/o according to <i>mode</i>
<i>stat(path, array)</i>	return information about <i>path</i>
<i>unlink(path)</i>	unlink <i>path</i>

is performed much less frequently than the latter. The contribution of PDS is a portable hierarchical directory system with absolutely no time impact on i/o efficiency.

2. USAGE

PDS is packaged as a set of RATFOR [2, 3] (and hence FORTRAN) functions and subroutines, which is loaded with the program or system that uses it. PDS is useful in programs or systems that use named files extensively, or where a machine-independent means of naming files and a flexible directory structure would enhance utility. Another use would be in a set of tools, such as those described in [3], implemented on computers with limited file systems.

The following program, *saveall*, illustrates a typical use of PDS. It copies all of the files in the current directory to files of the same name in the directory *../backup*. *saveall* is written in RATFOR in the style of [3].

```
# saveall - save all files in ../backup
character dline(MAXLINE), bname(MAXLINE)
integer fd, fdi, fdo
integer openf, creatf, getlin
string dot "."
string backup "../backup/"

fd = openf(dot, READ)
while (getlin(dline, fd) ~= EOF) {
    fdi = openf(dline(5), READ)
    call strcat(backup, dline(5), bname)
    fdo = creatf(bname, WRITE)
    call fcopy(fdi, fdo)
    call close(fdi)
    call close(fdo)
}
call close(fd)
end
```

Uppercase names denote defined constants; for example, a typical value for *MAXLINE* might be 80. The **string** statement declares a character array large enough to accommodate the indicated character string. The program begins by opening the current directory for reading using *openf*. The **while** loop reads the directory, line-by-line, until end-of-file, placing each line in the character array *dline*. As described in the next section, each line in a directory contains a file name beginning in column 5. Thus the body of the **while** loop opens each file for reading (*openf*), constructs the path name of the backup copy (*strcat*), creates the backup file and opens it for writing (*creatf*), copies the file (*fcopy*), and finally closes the original and backup files (*close*).

This program makes use of the i/o routines described in [3] (*getlin* and *close*), but any other set of routines—including standard FORTRAN **read** and **write** statements—could be used.

Another example of the use of PDS is as a simple command preprocessor that translates machine-independent system commands to the appropriate host command sequence. In commands, files are referred to by their PDS names and are

mapped to the corresponding host names during the generation of host commands. For example, the command

```
list files . . .
```

lists the named files on the line printer. The preprocessor translates the **list** into whatever is appropriate on the host system. For example,

```
list ../backup/saveall.r saveall.r
```

results in the DEC-10 command sequence

```
r queue
  ilist.120, save.rat
```

where “ilist.120” and “save.rat” are the host names for “../backup/saveall.r” and “saveall.r”, respectively.

3. IMPLEMENTATION

PDS is designed to use only simple sequential i/o (i.e., FORTRAN i/o). It is implemented in RATFOR for portability reasons and because FORTRAN is the only widely available high-level language to which calls from other languages can be made.

The implementation is similar to the implementation of the UNIX file system [7]. The basic technique is to separate the information about a file from the presence of that file in a directory. An *i-node* is associated with each file and directory, and it contains all of the information, including the host name, concerning the file. All of the i-nodes are stored in a single file in which i-node *n* is line *n*. This file, referred to as the *i-list*, is stored in character format so that it can be read easily with FORTRAN i/o or its equivalent.

The i-list maps i-node number, or i-number, to a host name. Directories provide a map of PDS name to i-number, thereby completing the mapping from PDS name to host name. A directory is a file of lines, each line containing an i-number and the corresponding PDS name.

Further implementation details are given in [1].

4. INSTALLATION

PDS consists of about 1000 lines of RATFOR including comments. Of this total, however, 300 lines are utility subroutines commonly available in a RATFOR environment. After preprocessing by RATFOR, the resulting FORTRAN code is about 1200 lines in length. The resulting FORTRAN conforms to the portable subset of ANSI standard FORTRAN as defined by PFORT [6].

On a DEC-10 (36-bit words, 512-word pages), the code area for the entire system occupies 6 pages and the data area occupies 17 pages. Small adjustments in the size of the data area can be made by changing various parameters, although the tendency seems to be to *increase* those parameters.

As mentioned in the previous section, PDS operates using sequential i/o only. It is therefore possible to use FORTRAN i/o, although in practice modifications are necessary to avoid having to use FORTRAN unit numbers and to make use of named host files. PDS is written to use the i/o interface described in [3]; a FORTRAN version of these routines is provided along with suggested modifications.

Installation of the system using the FORTRAN version of the i/o interface can be accomplished in 1–2 man-days. The implementation of more sophisticated i/o systems requires substantial time investment, typically 3–6 man-months depending on the target system. This is unnecessary unless heavy use of the RATFOR i/o interface is anticipated. System-specific i/o systems are supplied for the DEC-10 and Cyber 175.

PDS can be modified to use standard FORTRAN i/o facilities directly (i.e., **read** and **write** statements). The modification requires replacing parts of the RATFOR i/o interface—the 300 lines of code mentioned above—with the appro-

priate standard FORTRAN i/o statements. While these modifications have *not* been made, similar experience suggests that 1-2 man-weeks is a conservative estimate of the effort required.

PDS is distributed with the following components:

PDS written in RATFOR
PDS written in FORTRAN (RATFOR output)
RATFOR and RATFOR i/o system written in RATFOR
RATFOR and RATFOR i/o system written in FORTRAN
DEC-10 i/o system
Cyber 175 i/o system

REFERENCES

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2. KERNIGHAN, B.W. Ratfor—A preprocessor for a rational Fortran. *Softw. Pract. Exper.* 5, 4 (Dec. 1975), 396-406.
3. KERNIGHAN, B.W., AND PLAUGER, P.J. *Software Tools*. Addison-Wesley, Reading, Mass., 1976.
4. RITCHIE, D.M., AND THOMPSON, K. The UNIX timesharing system. *Commun. ACM* 17, 7 (July 1974), 365-375.
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6. RYDER, B.G. The PFORT verifier. *Softw. Pract. Exper.* 4, 4 (Dec. 1974), 359-377.
7. THOMPSON, K. UNIX implementation. *Bell Syst. Tech. J.* 57, 6 (July 1978), 1931-1946.

ALGORITHM

[The complete algorithm is available from the ACM Algorithms Distribution Service].

ALGORITHM 569

COLSYS: Collocation Software for Boundary-Value ODEs [D2]

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and

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Key Words and Phrases: ordinary differential equations, boundary-value problems, collocation, B-spline, mesh selection, error estimates, damped Newton's method, general-purpose code

CR Categories: 5.17

Language: FORTRAN

DESCRIPTION

This package is a complement to [1] where its usage is described and demonstrated.

COLSYS uses the following non-American National Standard FORTRAN conventions:

- (1) No run-time subscript range checking.
- (2) Simple expressions used as subscripts.
- (3) Specifications of array values in DATA statements (used in subroutine CONSTS) by referring to the array name only.

REFERENCES

1. ASCHER, U., CHRISTIANSEN, J., AND RUSSELL, R.D. Collocation software for boundary value ODEs. *ACM Trans. Math. Softw.* 7, 2 (June 1981), 209-222

ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```

SUBROUTINE COLSYS (NCOMP, M, ALEFT, ARIGHT, ZETA, IPAR, LTOL,
1      TOL, FIXPNT, ISPACE, FSPACE, IFLAG, FSUB,
2      DFSUB, GSUB, DGSUB, SOLUTN)

```

C

C

C*****

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```

C
C
C   PURPOSE
C
C   SUBROUTINE COLSYS SOLVES A MULTI-POINT BOUNDARY VALUE
C   PROBLEM FOR A MIXED ORDER SYSTEM OF ODE-S GIVEN BY
C
C       (M(I))
C       U      = F ( X; Z(U(X)) )      I = 1, ... ,NCOMP
C       I      I
C
C                               ALEFT .LT. X .LT. ARIGHT,
C
C
C       G ( ZETA(J); Z(U(ZETA(J))) )= 0  J = 1, ... ,MSTAR
C       J
C                               MSTAR=M(1)+M(2)+...+M(NCOMP),
C
C
C   WHERE
C       U = (U1, U2, ... ,UNCOMP)T IS THE EXACT SOLUTION VECTOR
C
C       (MI)
C       UI IS THE MI=M(I) TH DERIVATIVE OF UI
C
C       Z(U(X)) = (U1(1)(X), U1(M1-1)(X), ..., U1(M1-1)(X), ..., UNCOMP(MNCOMP-1)(X))T
C
C       F (X, Z(U)) IS A (GENERALLY) NONLINEAR FUNCTION OF
C       I
C       Z(U)=Z(U(X)).
C
C       G (ZETA(J); Z(U)) IS A (GENERALLY) NONLINEAR BOUNDARY
C       J
C       CONDITION.
C
C   THE BOUNDARY POINTS SATISFY
C       ALEFT .LE. ZETA(1) .LE. ... .LE. ZETA(MSTAR) .LE. ARIGHT
C
C   THE ORDERS MI OF THE DIFFERENTIAL EQUATIONS SATISFY
C       M1 .LE. M2 .LE. ... .LE. MNCOMP .LE. 4.
C
C*****
C
C   *****      INPUT TO COLSYS      *****
C
C   VARIABLES
C
C   NCOMP - NO. OF DIFFERENTIAL EQUATIONS (NCOMP .LE. 20)
C
C   M(J) - ORDER OF THE J-TH DIFFERENTIAL EQUATION ( M(J).LE.M(J+1)
C           AND MSTAR = M(1) + ... + M(NCOMP) .LE. 40 )
C
C   ALEFT - LEFT END OF INTERVAL
C
C   ARIGHT - RIGHT END OF INTERVAL
C
C   ZETA(J) - J-TH SIDE CONDITION POINT (BOUNDARY POINT). MUST
C             HAVE ZETA(J) .LE. ZETA(J+1)
C
C   IPAR - AN INTEGER ARRAY DIMENSIONED AT LEAST 11.
C          A LIST OF THE PARAMETERS IN IPAR AND THEIR MEANING FOLLOWS.
C          SOME PARAMETERS ARE RENAMED IN COLSYS, THESE NEW NAMES ARE
C          GIVEN IN PARENTHESES.
C
C   IPAR(1) ( = NONLIN )
C           = 0 IF THE PROBLEM IS LINEAR
C           = 1 IF THE PROBLEM IS NONLINEAR
C
C   IPAR(2) = NO. OF COLLOCATION POINTS PER SUBINTERVAL (= K )
C           WHERE M(NCOMP) .LT. K .LE. 7 . IF IPAR(2)=0 THEN
C           COLSYS SETS K = MAX ( M(NCOMP)+1, 5-M(NCOMP) )
C
C

```

```

C      IPAR(3) = NO. OF SUBINTERVALS IN THE INITIAL MESH ( = N ).
C          IF IPAR(3) = 0 THEN COLSYS ARBITRARILY SETS N = 5.
C
C      IPAR(4) = NO. OF SOLUTION AND DERIVATIVE TOLERANCES. ( = NTOL )
C          WE REQUIRE 0 .LT. NTOL .LE. MSTAR.
C
C      IPAR(5) = DIMENSION OF FSPACE. ( = NDMF )
C
C      IPAR(6) = DIMENSION OF ISPACE. ( = NDMI )
C
C      IPAR(7) - OUTPUT CONTROL ( = IPRINT )
C          = -1 FOR FULL DIAGNOSTIC PRINTOUT
C          = 0 FOR SELECTED PRINTOUT
C          = 1 FOR NO PRINTOUT
C
C      IPAR(8) ( = IREAD )
C          = 0 CAUSES COLSYS TO GENERATE A UNIFORM INITIAL MESH.
C          = 1 IF THE INITIAL MESH IS PROVIDED BY THE USER. IT
C            IS DEFINED IN FSPACE AS FOLLOWS: THE MESH
C            ALEFT=X(1).LT.X(2).LT. . . . .LT.X(N).LT.X(N+1)=ARIGHT
C            WILL OCCUPY FSPACE(1), . . . , FSPACE(N+1). THE
C            USER NEEDS TO SUPPLY ONLY THE INTERIOR MESH
C            POINTS FSPACE(J) = X(J), J = 2, . . . , N.
C          = 2 IF THE INITIAL MESH IS SUPPLIED BY THE USER
C            AS WITH IPAR(8)=1, AND IN ADDITION NO ADAPTIVE
C            MESH SELECTION IS TO BE DONE.
C
C      IPAR(9) ( = IGUESS )
C          = 0 IF NO INITIAL GUESS FOR THE SOLUTION IS
C            PROVIDED.
C          = 1 IF AN INITIAL GUESS IS PROVIDED BY THE USER
C            IN SUBROUTINE SOLUTN.
C          = 2 IF AN INITIAL MESH AND APPROXIMATE SOLUTION
C            COEFFICIENTS ARE PROVIDED BY THE USER IN FSPACE.
C            (THE FORMER AND NEW MESH ARE THE SAME).
C          = 3 IF A FORMER MESH AND AN APPROXIMATE SOLUTION
C            COEFFICIENTS ARE PROVIDED BY THE USER IN FSPACE,
C            AND THE NEW MESH IS TO BE TAKEN TWICE AS COARSE.
C          = 4 IF IN ADDITION TO A FORMER INITIAL MESH AND AN
C            APPROXIMATE SOLUTION COEFFICIENTS, A NEW MESH
C            IS PROVIDED IN FSPACE AS WELL.
C            (SEE DESCRIPTION OF OUTPUT FOR FURTHER DETAILS
C            ON IGUESS = 2, 3, AND 4.)
C
C      IPAR(10)= 0 IF THE PROBLEM IS REGULAR
C          = 1 IF THE FIRST RELAX FACTOR IS =RSTART, AND THE
C            NONLINEAR ITERATION DOES NOT RELY ON PAST COVERAGE
C            (USE FOR AN EXTRA SENSITIVE NONLINEAR PROBLEM ONLY).
C          = 2 IF WE ARE TO RETURN IMMEDIATELY UPON (A) TWO
C            SUCCESSIVE NONCONVERGENCES, OR (B) AFTER OBTAINING
C            ERROR ESTIMATE FOR THE FIRST TIME.
C
C      IPAR(11)= NO. OF FIXED POINTS IN THE MESH OTHER THAN
C          ALEFT AND ARIGHT. ( = NFXPNT , THE DIMENSION OF FIXPNT)
C
C      LTOL - AN ARRAY OF DIMENSION IPAR(4). LTOL(J) = L SPECIFIES
C            THAT THE J-TH TOLERANCE IN TOL CONTROLS THE ERROR
C            IN THE L-TH COMPONENT OF Z(U). ALSO REQUIRE THAT
C            1.LE.LTOL(1).LT.LTOL(2).LT. . . . .LT.LTOL(NTOL).LE.MSTAR
C
C      TOL - AN ARRAY OF DIMENSION IPAR(4). TOL(J) IS THE
C            ERROR TOLERANCE ON THE LTOL(J) -TH COMPONENT
C            OF Z(U). THUS, THE CODE ATTEMPTS TO SATISFY
C            FOR J=1, . . . , NTOL ON EACH SUBINTERVAL
C            ABS(Z(V)-Z(U)) .LE. TOL(J)*Z(U) +TOL(J)
C            LTOL(J) LTOL(J)
C            IF V(X) IS THE APPROXIMATE SOLUTION VECTOR.
C
C      FIXPNT - AN ARRAY OF DIMENSION IPAR(11). IT CONTAINS
C            THE POINTS, OTHER THAN ALEFT AND ARIGHT, WHICH
C            ARE TO BE INCLUDED IN EVERY MESH.
C
C      ISPACE - AN INTEGER WORK ARRAY OF DIMENSION IPAR(6).
C            ITS SIZE PROVIDES A CONSTRAINT ON NMAX,
C            THE MAXIMUM NUMBER OF SUBINTERVALS. CHOOSE
C            IPAR(6) ACCORDING TO THE FORMULA

```

```

C          IPAR(6) .GE. NMAX*NSIZEI
C          WHERE
C          NSIZEI = 3 + KDM - NREC
C          WITH
C          KDM = KD + MSTAR ; KD = K * NCOMP ;
C          NREC = NO. OF RIGHT END BOUNDARY CONDITIONS.
C
C          FSPACE - A REAL WORK ARRAY OF DIMENSION IPAR(5).
C          ITS SIZE PROVIDES A CONSTRAINT ON NMAX.
C          CHOOSE IPAR(5) ACCORDING TO THE FORMULA
C          IPAR(5) .GE. NMAX*NSIZEF
C          WHERE
C          NSIZEF = 4 + K + 2 * KD + (4+2*K) * MSTAR +
C                  (KDM-NREC) * (KDM+1).
C
C          IFLAG - THE MODE OF RETURN FROM COLSYS.
C          = 1 FOR NORMAL RETURN
C          = 0 IF THE COLLOCATION MATRIX IS SINGULAR.
C          =-1 IF THE EXPECTED NO. OF SUBINTERVALS EXCEEDS STORAGE
C             SPECIFICATIONS.
C          =-2 IF THE NONLINEAR ITERATION HAS NOT CONVERGED.
C          =-3 IF THERE IS AN INPUT DATA ERROR.
C
C*****
C          *****      USER SUPPLIED EXTERNAL SUBROUTINES      *****
C
C          FSUB - NAME OF SUBROUTINE FOR EVALUATING  $F(X, Z(U(X))) =$ 
C                 $(F_1, \dots, F_{NCOMP})$  AT A POINT X IN (ALEFT, ARIGHT). IT
C                SHOULD HAVE THE HEADING
C
C                SUBROUTINE FSUB ( X , Z , F )
C
C                WHERE F IS THE VECTOR CONTAINING THE VALUE OF  $F_i(X, Z(U))$ 
C                IN THE I-TH COMPONENT AND  $Z(U(X)) = (Z(1), \dots, Z(MSTAR))$ 
C                IS DEFINED AS ABOVE UNDER PURPOSE .
C
C          DFSUB - NAME OF SUBROUTINE FOR EVALUATING THE JACOBIAN OF
C           $F(X, Z(U))$  AT A POINT X. IT SHOULD HAVE THE HEADING
C
C                SUBROUTINE DFSUB ( X , Z , DF )
C
C                WHERE  $Z(U(X))$  IS DEFINED AS FOR FSUB AND THE (NCOMP) BY
C                (MSTAR) ARRAY DF SHOULD BE FILLED BY THE PARTIAL DERIV-
C                ATIVES OF F, VIZ, FOR A PARTICULAR CALL ONE CALCULATES
C                 $DF(I, J) = DF_i / DZ_j$ , I=1, ..., NCOMP
C                J=1, ..., MSTAR.
C
C          GSUB - NAME OF SUBROUTINE FOR EVALUATING THE I-TH COMPONENT OF
C           $G(X, Z(U(X))) = G_I(ZETA(I), Z(U(ZETA(I))))$  AT A POINT X =
C                ZETA(I) WHERE 1. LE. I. LE. MSTAR. IT SHOULD HAVE THE HEADING
C
C                SUBROUTINE GSUB ( I , Z , G )
C
C                WHERE Z(U) IS AS FOR FSUB, AND I AND  $G = G_I$  ARE AS ABOVE.
C                NOTE THAT IN CONTRAST TO F IN FSUB, HERE
C                ONLY ONE VALUE PER CALL IS RETURNED IN G.
C
C          DGSUB - NAME OF SUBROUTINE FOR EVALUATING THE I-TH ROW OF
C          THE JACOBIAN OF  $G(X, U(X))$ . IT SHOULD HAVE THE HEADING
C
C                SUBROUTINE DGSUB ( I , Z , DG )
C
C                WHERE Z(U) IS AS FOR FSUB, I AS FOR GSUB AND THE MSTAR-
C                VECTOR DG SHOULD BE FILLED WITH THE PARTIAL DERIVATIVES

```



```

C          OF G, VIZ, FOR A PARTICULAR CALL ONE CALCULATES
C          DG(I,J) = DGI / DZJ      J=1,...,MSTAR.
C
C
C          SOLUTN- NAME OF SUBROUTINE TO EVALUATE THE INITIAL
C          APPROXIMATION FOR Z(U(X)) AND FOR DMVAL(U(X))= VECTOR
C          OF THE MJ-TH DERIVATIVES OF U(X). IT SHOULD HAVE THE
C          HEADING
C
C                  SUBROUTINE SOLUTN ( X , Z , DMVAL )
C
C          NOTE THAT THIS SUBROUTINE IS NEEDED ONLY IF USING
C          IPAR(9) = 1, AND THEN ALL MSTAR COMPONENTS OF Z
C          AND NCOMP COMPONENTS OF DMVAL SHOULD BE SPECIFIED
C          FOR ANY X, ALEFT.LE. X.LE. ARIGHT .
C
C*****
C          *****          OUTPUT FROM COLSYS          *****
C
C          UPON RETURN FROM COLSYS , THE USER MAY PRODUCE THE
C          SOLUTION VECTOR Z( U(X) ) AT A POINT X, ALEFT.LE.X.LE.ARIGHT
C          BY CALLING :
C
C                  CALL APPSLN ( X, Z, FSPACE, ISPACE )
C
C          THIS SETS UP A STANDARD CALL TO APPROX . FOR A MORE
C          EFFICIENT OR SOPHISTICATED RETRIEVAL OF THE SOLUTION
C          VALUES, CALL APPROX DIRECTLY (SEE DOCUMENTATION IN
C          APPROX - THE PARAMETERS NEEDED IN THE CALL TO APPROX
C          BY THE USER ARE SAVED IN ISPACE AND FSPACE BEFORE
C          COLSYS RETURNS).
C
C          IN ORDER TO SAVE THE COEFFICIENTS OF THE SOLUTION FOR LATER
C          REFERENCE,      ISPACE(1), ... , ISPACE(7+MSTAR) AND
C          FSPACE(1), ... , FSPACE(ISPACE(7)) SHOULD BE
C          SAVED, SINCE THESE ARE USED IN THE CALL TO APPSLN (APPROX).
C
C          ONE CAN ALSO USE THE FORMERLY OBTAINED APPROXIMATE
C          SOLUTION AS A FIRST APPROXIMATION FOR THE NONLINEAR ITERATION
C          ON A NEW PROBLEM (E.G. FOR CONTINUATION PURPOSES). THIS
C          INVOLVES USING Iguess = 2, 3, OR 4, AS FOLLOWS:
C
C          FOR Iguess= 2 OR 3, THE USER SHOULD PUT THE ABOVE SAVED
C          VALUES BACK INTO FSPACE(1),... ,FSPACE(ISPACE(6)).
C          THE SIZE OF THE FORMER MESH, NOLD, IS PROVIDED IN IPAR(3). IF
C          Iguess=2 THEN THE SIZE OF THE NEW MESH, N, IS TAKEN TO BE =NOLD.
C          IF Iguess=3 THEN N := NOLD/2 AND THE NEW MESH IS TO BE TWICE AS
C          COARSE.
C          FOR Iguess=4, PUT N IN IPAR(3) AND NOLD IN ISPACE(1). THE
C          VALUES OF THE FORMER SOLUTION, SAVED AS DESCRIBED ABOVE,
C          SHOULD BE PUT INTO FSPACE(N+2),... ,FSPACE(ISPACE(6)+N+1), AND
C          A NEW MESH UNRELATED TO THE FORMER ONE IS PRESCRIBED IN
C          FSPACE(1),... ,FSPACE(N+1).
C
C*****
C          *****          PACKAGE SUBROUTINES          *****
C
C          THE FOLLOWING DESCRIPTION GIVES A BRIEF OVERVIEW OF HOW THE
C          PROCEDURE IS BROKEN DOWN INTO THE SUBROUTINES WHICH MAKE UP
C          THE PACKAGE CALLED COLSYS . FOR FURTHER DETAILS THE
C          USER SHOULD REFER TO DOCUMENTATION IN THE VARIOUS SUBROUTINES
C          AND TO THE REFERENCES CITED ABOVE.
C
C          THE SUBROUTINES FALL INTO FOUR GROUPS:
C
C          PART 1 - THE MAIN STORAGE ALLOCATION AND PROGRAM CONTROL SUBROUTINES.
C
C          COLSYS - TESTS INPUT VALUES, DOES INITIALIZATION AND BREAKS UP
C          THE WORK AREAS, FSPACE AND ISPACE, INTO THE ARRAYS
C          USED BY THE PROGRAM.
C

```

```
C   CONTRL - IS THE ACTUAL DRIVER OF THE PACKAGE. THIS ROUTINE
C           CONTAINS THE STRATEGY FOR NONLINEAR PROBLEMS.
C
C PART 2 - MESH SELECTION AND ERROR ESTIMATION SUBROUTINES
C
C   CONSTS - IS CALLED ONCE BY COLSYS TO INITIALIZE CONSTANTS
C           WHICH ARE USED FOR ERROR ESTIMATION AND MESH SELECTION.
C
C   NEWMSH - GENERATES MESHES. IT CONTAINS THE TEST TO DECIDE
C           WHETHER OR NOT TO REDISTRIBUTE A MESH.
C
C   ERRCHK - PRODUCES ERROR ESTIMATES AND CHECKS AGAINST THE
C           TOLERANCES AT EACH SUBINTERVAL
C
C PART 3 - COLLOCATION SYSTEM SET-UP SUBROUTINES
C
C   LSYSLV - CONTROLS THE SET-UP AND SOLUTION OF THE LINEAR
C           ALGEBRAIC SYSTEMS OF COLLOCATION EQUATIONS WHICH
C           ARISE AT EACH NEWTON ITERATION.
C
C   BLDBLK - IS USED BY LSYSLV TO SET UP THE EQUATION(S) ASSOCIATED
C           WITH A SIDE CONDITION POINT OR A COLLOCATION POINT.
C
C PART 4 - B-SPLINE SUBROUTINES
C
C   APPSLN - SETS UP A STANDARD CALL TO APPROX
C
C   APPROX - EVALUATES A PIECEWISE POLYNOMIAL SOLUTION.
C
C   BSPFIX - EVALUATES THE MESH INDEPENDENT B-SPLINES
C           (I. E. THE FIXED B-SPLINES)
C
C   BSPVAR - EVALUATES THE MESH DEPENDENT B-SPLINES (I. E. THE
C           VARYING B-SPLINES)
C
C   BSPDER - GENERATES VALUES FOR THE DERIVATIVES NEEDED TO SET
C           UP THE COLLOCATION EQUATIONS.
C
C   APPDIF - GENERATES A DIVIDED DIFFERENCE TABLE FROM THE B-SPLINE
C           COEFFICIENTS FOR A COLLOCATION SOLUTION. THE TABLE
C           IS USED IN APPROX
C
C   HORDER - EVALUATES THE HIGHEST ORDER DERIVATIVES OF THE
C           CURRENT COLLOCATION SOLUTION USED FOR MESH REFINEMENT.
C
C TO SOLVE THE LINEAR SYSTEMS OF COLLOCATION EQUATIONS
C CONSTRUCTED IN PART 3, COLSYS USES THE PACKAGE
C SOLVEBLOK OF DE BOOR - WEISS (TO APPEAR IN TOMS).
```

ALGORITHM 570

LOPSI: A Simultaneous Iteration Algorithm for Real Matrices [F2]

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and

ALAN JENNINGS

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Key Words and Phrases: eigenvalues, eigenvectors, simultaneous iteration, real unsymmetric matrices, large sparse matrices

CR Categories: 5.14

Language: FORTRAN

DESCRIPTION

The algorithm given here is a complement to [1] where the description, test results, and references are given.

REFERENCES

1. STEWART, W.J., AND JENNINGS, A. A simultaneous iteration algorithm for real matrices. *ACM Trans. Math. Softw.* 7, 2 (June 1981), 184-198.

ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

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SUBROUTINE LOPSI(N, A, IA, IROW, ICOL, LR, IR, M, IS, TOL, ITMAX, LMAX, U, IU,
1          V, W, G, H, B, P, IQ, INT, X, Y, ERR, ITTOT, INACT, IFAIL)
C
C
C
C      ** LOPSI IS A SIMULTANEOUS ITERATION ALGORITHM WHICH
C      DETERMINES APPROXIMATIONS TO RIGHT OR LEFT EIGEN-
C      VECTORS CORRESPONDING TO THE DOMINANT SET OF
C      EIGENVALUES OF A REAL UNSYMMETRIC MATRIX A.      **
C
C          BY WILLIAM J. STEWART AND ALAN JENNINGS.
C
C
C
C      FORMAL PARAMETER LIST.

```

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```

C *****
C
C
C INPUT PARAMETERS (MUST BE SET BEFORE ENTRY AND ARE UNALTERED
C ***** BY THE EXECUTION OF THE ALGORITHM)
C
C
C N AN INTEGER QUANTITY, THE ORDER OF THE MATRIX A FOR WHICH
C THE PARTIAL EIGENSOLUTION IS TO BE OBTAINED.
C
C A A ONE-DIMENSIONAL REAL ARRAY OF LENGTH AT LEAST EQUAL
C TO IA. IT CONTAINS IN ARBITRARY ORDER THE NON-ZERO
C ELEMENTS OF THE MATRIX A.
C
C IA AN INTEGER QUANTITY, THE NUMBER OF NON-ZERO ELEMENTS IN
C THE MATRIX A.
C
C IROW A ONE-DIMENSIONAL INTEGER ARRAY OF LENGTH AT LEAST EQUAL TO
C IA. THE I-TH COMPONENT DENOTES THE ROW POSITION OF THE
C NON-ZERO ELEMENT STORED IN POSITION I OF ARRAY A.
C
C ICOL A ONE-DIMENSIONAL INTEGER ARRAY OF LENGTH AT LEAST EQUAL TO
C IA. THE I-TH COMPONENT DENOTES THE COLUMN POSITION OF THE
C NON-ZERO ELEMENT STORED IN POSITION I OF ARRAY A.
C
C LR AN INTEGER QUANTITY. IF LR=1 ON ENTRY, THEN THE RIGHT EIGEN-
C VECTORS CORRESPONDING TO THE DOMINANT EIGENVALUES WILL BE
C OBTAINED. OTHERWISE THE LEFT SET WILL BE OBTAINED.
C
C IR AN INTEGER QUANTITY, THE NO. OF VECTORS REQUIRED ACCURATELY.
C
C M AN INTEGER QUANTITY, THE NO. OF TRIAL VECTORS EMPLOYED.
C M SHOULD NORMALLY BE CHOSEN SO THAT M >= IR+2
C
C IS AN INTEGER QUANTITY, THE NUMBER OF TRIAL VECTORS FOR WHICH
C INITIAL APPROXIMATIONS ARE SUPPLIED BY THE USER.
C
C TOL A REAL QUANTITY, THE TOLERANCE DEMANDED OF THE EIGENVECTORS
C CORRESPONDING TO THE IR DOMINANT EIGENVALUES.
C THE EIGENVALUES ARE NORMALIZED TO HAVE LARGEST COMPONENT
C UNITY AND "LOPSI" AIMS FOR EACH COMPONENT TO BE IN ERROR
C BY LESS THAN TOL. IN GENERAL IT CAN BE EXPECTED THAT
C THE CORRESPONDING EIGENVALUES HAVE A RELATIVE ERROR
C LESS THAN TOL.
C
C ITMAX AN INTEGER QUANTITY, AN UPPER BOUND ON THE NUMBER OF MATRIX
C BY VECTOR MULTIPLICATIONS TO BE PERFORMED
C
C LMAX AN INTEGER QUANTITY, THE MAXIMUM NUMBER OF PREMULTIPLIC-
C ATIONS PERMITTED BETWEEN TWO REORIENTATIONS
C THE VALUE LMAX = 10 SHOULD BE ADEQUATE FOR MOST PROBLEMS
C HOWEVER SET LMAX = 1 WHEN EXPECTING FAST CONVERGENCE
C FOR FURTHER INFORMATION SEE SECTION 3 OF THE PAPER
C
C IU AN INTEGER QUANTITY, THE DECLARED FIRST DIMENSION OF
C ARRAYS U, V AND W.
C
C IG AN INTEGER QUANTITY, THE DECLARED FIRST DIMENSION
C OF WORK ARRAYS G, H, B, P.
C
C INPUT AND OUTPUT PARAMETER
C *****
C
C U A TWO-DIMENSIONAL REAL ARRAY OF AT LEAST (N, M) ELEMENTS. ON
C ENTRY IT CONTAINS USER SUPPLIED APPROXIMATIONS TO THE EIGEN-
C VECTORS CORRESPONDING TO THE IS DOMINANT EIGENVALUES.
C ON EXIT IT CONTAINS THE FINAL ESTIMATES TO THE EIGENVECTORS
C CORRESPONDING TO THE M DOMINANT EIGENVALUES.
C
C OUTPUT PARAMETERS
C *****
C
C

```

```

C      X, Y  ONE-DIMENSIONAL REAL ARRAYS OF LENGTH AT LEAST EQUAL TO M.
C           ON EXIT THE REAL AND IMAGINARY PARTS OF THE J-TH EIGENVALUE
C           PREDICTION ARE STORED IN X(J) AND Y(J) RESPECTIVELY.
C
C      ERR   A ONE-DIMENSIONAL REAL ARRAY OF AT LEAST M ELEMENTS. ON EXIT
C           ERR(J) DENOTES THE ESTIMATED ACCURACY OF THE J-TH VECTOR.
C           ERROR ESTIMATES, AS MEASURED BY THE MAXIMUM OF THE ABSOLUTE
C           DIFFERENCES IN THE COMPONENTS OF SUCCESSIVE APPROXIMATIONS,
C           ARE GIVEN FOR ALL M VECTORS INCLUDING THOSE THAT HAVE NOT
C           SATISFIED THE CONVERGENCE CRITERIA SPECIFIED UNDER TOL.
C
C      ITTOT AN INTEGER QUANTITY WHICH ON EXIT DENOTES THE
C           NUMBER OF PREMULIPLICATIONS PERFORMED.
C
C      INACT AN INTEGER QUANTITY WHICH ON EXIT DENOTES THE
C           NUMBER OF INTERACTION ANALYSES PERFORMED.
C
C      IFAIL AN INTEGER QUANTITY. ON EXIT IFAIL=0 IF FIRST IR VECTORS
C           HAVE SATISFIED THE CONVERGENCE CRITERIA, OTHERWISE IFAIL=1.
C
C
C      WORK ARRAYS
C      *****
C
C      V,W   TWO-DIMENSIONAL REAL WORK ARRAYS OF AT LEAST (N, M) ELEMENTS
C
C      G,H,  FOUR TWO-DIMENSIONAL REAL WORK ARRAYS
C      B,P   OF AT LEAST (M, M) ELEMENTS EACH.
C
C      INT   A ONE-DIMENSIONAL INTEGER WORK ARRAY OF AT LEAST M ELEMENTS.
C
C
C      DOUBLE PRECISION A, B, DIFF1, DIFF2, DLAM, DLAMM, DLAMR, DLAM1, EMAX
C      DOUBLE PRECISION ERR, FMAX, G, H, P, SIGMA, SUM, TOIT, TOL, TOT, TOT1, TOT2
C      DOUBLE PRECISION U, V, W, X, Y, EPMACH, EPS
C      DIMENSION A(IA), IROW(IA), ICOL(IA), U(IU, M), V(IU, M), W(IU, M)
C      DIMENSION G(IG, M), H(IG, M), B(IG, M), P(IG, M)
C      DIMENSION X(M), Y(M), INT(M), ERR(M)

```


ALGORITHM 571

Statistics for von Mises' and Fisher's Distributions of Directions: $I_1(x)/I_0(x)$, $I_{1.5}(x)/I_{0.5}(x)$, and Their Inverses [S14]

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Key Words and Phrases: direction statistics, von Mises distribution, Fisher distribution, modified Bessel function ratio, continued fraction, function inversion, Newton-Raphson
CR Categories: 5.5, 5.12
Language: FORTRAN

DESCRIPTION

Two pairs of FORTRAN functions provide for evaluation of the ratio of modified Bessel functions of the first kind, $A(\kappa) = I_1(\kappa)/I_0(\kappa)$ or $B(\kappa) = I_{1.5}(\kappa)/I_{0.5}(\kappa) = \coth \kappa - 1/\kappa$ or their inverses, $A^{-1}(R)$, $B^{-1}(R)$. $A(\kappa)$ is the expected value of the mean modulus \bar{R} of random unit vectors in two dimensions, distributed with concentration parameter κ in von Mises' [5] distribution of directions on a circle, whereas $B(\kappa) = E(\bar{R})$ for Fisher's [1] distribution of three-dimensional directions on a sphere. The inverse functions, $A^{-1}(\bar{R})$ and $B^{-1}(\bar{R})$ are corresponding maximum likelihood estimators of κ for observed \bar{R} .

Other versions of each function are fairly readily achieved, which provide some choice of precision level in the real valued result, summarized in Table I.

The function BESRAT is illustrated as a 9.3S version, which selects between continued fractions detailed in [3] for backward recursive evaluation from a depth determined by parameter values C1 and C2, assigned with CX to DATA constants as indicated in comment cards. Within the tested range of precision the values specified for these parameters prevent occurrence of zero divisors in evaluation of the continued fractions. A version with S decimal digit precision assignable-at-call (up to limitations imposed by processor precision) can be obtained by changing the first card to read FUNCTION BESRAT(V, S), by deleting C from column 1 of cards defining C1, C2, and CX, and by suppressing the DATA card.

The function VKAPPA achieves 8S results by means of different initial approximations and up to two Newton-Raphson improvements adapted for several subintervals of the range of valid arguments. The Newton-Raphson steps call on BESRAT in a 9.3S version to compensate for some loss of precision for \bar{R} near 1, such as occurs in Algorithm AS81 [6]. None of the five expressions

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Table I

FORTRAN Function	Argument ^a	Result	Precision ^b
BESRAT(V)	$0 \leq \kappa = V $	$A(\kappa)$ in (0, 1)	9.3 (5-14)
VKAPPA(R)	$0 \leq R < 1$	$A^{-1}(R) \geq 0$	8 (5)
SPHERR(CAPPA)	$0 \leq \kappa = \text{CAPPA} $	$B(\kappa)$ in (0, 1)	13 (5-20)
CAPPA3(R)	$0 \leq R < 1$	$B^{-1}(R) \geq 0$	16 (8, 25)

^a Definition of κ as the modulus of the actual parameter value avoids an error condition, whereas for invalid argument $R < 0$ or $R \geq 1$ the value -1.0 is returned as an error signal result.

^b Precision is quoted as the least number of correct significant decimal digits in test results on a CDC 7600 of the FORTRAN version illustrated (range of precision tested is noted in parentheses). It is assumed that argument values are represented exactly and that the result precision is not appreciably limited by that of the processor.

occurring as divisors in the code can become zero for arguments within the relevant subintervals, so that processor arithmetic overflow cannot occur.

For analysis of observational data, less precise versions may be adequate [2], such as a fast 5S version of VKAPPA, obtained by assigning the value 0.74 to V1 and 0.89 to V2 in the DATA statement and by replacing the two conditional statements before label 30 by columns 29 to 72 of either to leave one unconditional Newton-Raphson call of a 6S version of BESRAT. In any case, the backward recursive evaluation of the continued fractions reduces accumulation of roundoff error and enables a substantially faster and more compact code than previously published procedures [2, 6], which involve some ten or more evaluations of numerator and denominator of the Bessel function ratio.

The function SPHERR determines $B(\kappa)$ either by backward recursive evaluation of the continued fraction form of $\coth \kappa - 1/\kappa$ or, if $\kappa > 1$, by means of the exponential form of \coth , as outlined in [3]. Unnecessary evaluation of the exponential function and possible overflow for large argument is avoided by evaluating $B(\kappa)$ as $1 - 1/\kappa$ for large values of $\kappa > \text{BIGX} = 1.1513 \times S + 0.4$, which ensures S decimal digit precision in the result. The depth of recursion is also simply determined as a function of desired precision S , as indicated in comment cards defining DATA constants. A version with precision assignable-at-call may be obtained by including S as a formal parameter and replacing the DATA statement by the assignments $\text{SON4} = S/4$ and $\text{BIGX} = 1.1513 \times S + 0.4$.

The function CAPPA3 achieves 8S precision in the result $B^{-1}(R)$ by means of inverses of the continued fraction and exponential forms of $\coth \kappa - 1/\kappa$, empirically extended over different subintervals of the argument [3]. Unnecessary evaluation of the exponential function and possible overflow for argument R near 1.0 are avoided by accepting $1/(1 - R)$ as the result, if that value exceeds $\text{BIGX} = 1.1513 \times S + 0.4$. Of course, the effective precision of the argument and therefore that of the result are limited by that of the difference, $1 - R$.

Precision of the initial approximation may be extended by Newton-Raphson inversion of $B(\kappa)$, using SPHERR(-) to provide the Bessel function ratio to a precision of $S + 1$ decimals to allow for loss of up to one decimal digit precision in the improvement. A version with precision up to 8S is obtained by deleting statements between label 20 and 30 of the version below, which caters for precision up to 16S. Precision could be extended to 25S by repeating the statement preceding label 30 or 32S by repeating the three statements preceding label 30, changing the DATA constants W1, W2 from 10.0, 0.01 to 20.0 and 0.001, respectively, and adjusting the precision of the auxiliary procedure SPHERR(-) accordingly.

Test programs on a CDC 7600 checked precision claims using single-precision (48B) versions of BESRAT and VKAPPA and double-precision (96B) versions of SPHERR and CAPPA3. Tabulated values checked against tables to 5D [5] identified some errors in published tables [4], but otherwise confirmed the validity of the functions BESRAT, VKAPPA, SPHERR, and CAPPA3.

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ALGORITHM

[The complete listing is printed here and is available from the ACM Algorithms Distribution Service.]

```

FUNCTION BESRAT(V)
-----
C RETURNS BESRAT = A(K) FOR K = ABS(V), WHERE A(K) IS THE EXPECTED
C MODULUS OF THE MEAN VECTOR SUM OF UNIT VECTORS SAMPLED FROM THE
C VON MISES DISTRIBUTION OF DIRECTIONS IN 2D WITH PARAMETER = K.
C A(V) = THE RATIO OF MODIFIED BESSEL FUNCTIONS OF THE FIRST KIND
C OF ORDERS 1 AND 0, I. E., A(V) = I1(V)/I0(V).
-----
C
C ADJUST TO S DECIMAL DIGIT PRECISION BY SETTING DATA CONSTANTS -
C C1 = (S+9.0-8.0/S)*0.0351
C C2 = ((S-5.0)**3/180.0+S-5.0)/10.0
C CX = S*0.5 + 11.0
C FOR S IN RANGE (5,14). THUS FOR S = 9.3 :
C DATA C1/0.613/, C2/0.475/, CX/15.65/
C
C Y = 0.0
C X = ABS(V)
C IF (X.GT.CX) GO TO 20
C
C FOR SMALL X, RATIO = X/(2+X*X/(4+X*X/(6+X*X/(8+ ... )))
C N = INT((X+16.0-16.0/(X+C1+0.75))*C1)
C X = X*0.5
C XX = X*X
C DO 10 J = 1, N
10 Y = XX/(FLOAT(N-J+2)+Y)
C BESRAT = X/(1.0+Y)
C RETURN
C
C FOR LARGE X, RATIO = 1-2/(4X-1-1/(4X/3-2-1/(4X/5-2- ... )))
20 N = INT((68.0/X+1.0)*C2) + 1
C X = X*4.0
C XX = FLOAT(N*2+1)
C DO 30 J = 1, N
C Y = XX/((-2.0-Y)*XX+X)
30 XX = XX - 2.0
C BESRAT = 1.0-2.0/(X-1.0-Y)
C RETURN
C END

FUNCTION VKAPPA(R)
-----
C RETURNS VKAPPA = THE MAXIMUM LIKELIHOOD ESTIMATE OF 'KAPPA', THE
C CONCENTRATION PARAMETER OF VON MISES' DISTRIBUTION OF DIRECTIONS
C IN 2 DIMENSIONS, CORRESPONDING TO A SAMPLE MEAN VECTOR MODULUS R.
C VKAPPA = K(A), THE INVERSE FUNCTION OF A(K) = RATIO OF MODIFIED
C BESSEL FUNCTIONS OF THE FIRST KIND, VIZ., A(K) = I1(K)/I0(K).
-----
C
C FOR 8S (SIGNIFICANT DECIMAL DIGITS) PRECISION AUXILIARY ROUTINE
C FUNCTION BESRAT(V) MUST BE SET TO AT LEAST 9.3S
C DATA V1/0.642/, V2/0.95/
C A = R
C S = -1.0
C

```

```

C ERROR SIGNAL: VALUE -1.0 RETURNED IF ARGUMENT -VE OR 1.0 OR MORE.
  IF (A.LT.0.0 .OR. A.GE.1.0) GO TO 30
  Y = 2.0/(1.0-A)
  IF (A.GT.0.85) GO TO 10
C
C FOR R BELOW 0.85 USE ADJUSTED INVERSE TAYLOR SERIES.
  X = A*A
  S = (((A-5.6076)*A+5.0797)*A-4.6494)*Y*X - 1.0
  S = (((S*X+15.0)*X+60.0)*X/360.0+1.0)*X-2.0)*A/(X-1.0)
  IF (V1-A) 20,20,30
C
C FOR R ABOVE 0.85 USE CONTINUED FRACTION APPROXIMATION.
10 IF (A.GE.0.95) X = 32.0/(120.0*A-131.5+Y)
  IF (A.LT.0.95) X = (-2326.0*A+4317.5526)*A-2001.035224
  S = (Y+1.0+3.0/(Y-5.0-12.0/(Y-10.0-X)))*0.25
  IF (A.GE.V2) GO TO 30
C
C FOR R IN (0.642,0.95) APPLY NEWTON-RAPHSON. TWICE IF R IN
C (0.75,0.875), FOR 85 PRECISION, USING APPROXIMATE DERIVATIVE -
20 Y = ((0.00048*Y-0.1589)*Y+0.744)*Y - 4.2932
  IF (A.LE.0.875) S = (BESRAT(S)-A)*Y + S
  IF (A.GE.0.75) S = (BESRAT(S)-A)*Y + S
30 VKAPPA = S
  RETURN
  END

```

FUNCTION SPHERR(CAPPA)

```

C -----
C RETURNS SPHERR = B(K) FOR K = ABS(CAPPA). B(K) IS THE EXPECTED
C MODULUS OF THE MEAN VECTOR SUM OF UNIT VECTORS SAMPLED FROM THE
C FISHER DISTRIBUTION OF DIRECTIONS IN 3D WITH PARAMETER = CAPPA.
C B(K) = THE RATIO OF MODIFIED BESSEL FUNCTIONS OF THE FIRST KIND
C OF ORDERS 3/2 AND 1/2, EQUIVALENT TO COTH(K) - 1/K.
C -----
C
C FOR S DECIMAL DIGIT PRECISION AND TO AVOID EXPONENTIAL OVERFLOW
C SET SON4 = S/4 AND BIGX = 1.1513*S+0.4, E.G., FOR 488 = 14.45S;
  DATA SON4/3.61/, BIGX/17.0/
C
  T = 0.0
  X = ABS(CAPPA)
  IF (X.GT.1.0) GO TO 20
C
C FOR SMALL X EVALUATE GAUSS CONTINUED FRACTION X/(3+X*X/(5+...))
C UP FROM J-TH LEVEL, WHERE N=2*J+1 YIELDS S DECIMALS PRECISION.
  N = INT((X+0.88)*SON4)*2 + 5
  XX = X*X
10 T = XX/(FLOAT(N)+T)
  N = N - 2
  IF (N.GT.3) GO TO 10
  SPHERR = X/(T+3.0)
  RETURN
C
C FOR LARGE X USE EXPONENTIAL FORM OF COTH(X) - 1/X
20 IF (X.LT.BIGX) T = 2.0/(EXP(2.0*X)-1.0)
  SPHERR = T + 1.0 - 1.0/X
  RETURN
  END

```

FUNCTION CAPPA3(R)

```

C -----
C RETURNS CAPPA3 = THE MAXIMUM LIKELIHOOD ESTIMATE OF 'KAPPA', THE
C CONCENTRATION PARAMETER OF THE FISHER DISTRIBUTION OF DIRECTIONS
C IN 3 DIMENSIONS, CORRESPONDING TO A SAMPLE MEAN VECTOR MODULUS R.
C CAPPA3 = K(B), THE INVERSE FUNCTION OF B(K) = RATIO OF MODIFIED
C BESSEL FUNCTIONS OF THE FIRST KIND OF ORDERS 3/2 AND 1/2.
C -----
C
C FOR PRECISION UP TO 85 (SIGNIFICANT DECIMAL DIGITS) OMIT THREE
C LINES FOLLOWING STATEMENT LABELED 20. FOR GREATER PRECISION UP
C TO 16S THE AUXILIARY SUBPROGRAM, FUNCTION SPHERR(X), IS NEEDED
C WITH PRECISION ONE DECIMAL DIGIT GREATER THAN FOR CAPPA3(R).
C TO AVOID EXPONENTIAL OVERFLOW SET BIGX = 1.1513*S + 0.4; FOR 488
  DATA BIGX /17.0/, W1/10.0/, W2/0.01/
  Y = R

```

```

      X = -1.0
C
C ERROR SIGNAL: VALUE -1.0 RETURNED IF ARGUMENT -VE OR 1.0 OR MORE
      IF (Y.LT.0.0 .OR. Y.GE.1.0) GO TO 30
      IF (Y.LT.0.5) GO TO 10
C
C FOR LARGE R APPROX. INVERSE OF R = COTH(K) - 1/K. (YIELDS 8.1S)
      X = 1.0/(1.0-Y)
      IF (X.GT.BIGX) GO TO 30
      S = 2.0*X
      T = EXP(S) - S*S - 1.0
      IF (Y.LT.0.8) T = (((0.00254*S-0.071042)*S+0.6943388)*S
                        -2.3816184)*S + 0.1508478 - 0.14789*Y + T
      X = T*X/(T+S)
      IF (X-W1) 20,20,30
C
C FOR SMALL R USE INVERSE GAUSS CONTINUED FRACTION. (YIELDS 8.4S)
10  X = 3.0*Y
      S = X*X
      T = 12.375
      IF (X.GT.0.7) T = ((5.0*X-14.74)*X+16.5198)*X+6.2762
      X = X/(1.0-S/(15.0-4.0*S/(7.0-S/(5.0-S/T))))
C
C ONE STAGE NEWTON-RAPHSON INVERSION DOUBLES PRECISION.
20  IF (X.LT.W2) GO TO 30
      S = EXP(X)
      S = S*2.0/(S*S-1.0)
      X = X + (Y-SPHERR(X))/(1.0/(X*X)-S*S)
C
30  CAPPAS = X
      RETURN
      END

```


ALGORITHM 572

Solution of the Helmholtz Equation for the Dirichlet Problem on General Bounded Three-Dimensional Regions [D3]

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and
OLOF WIDLUND
New York University

Key Words and Phrases: Helmholtz equation, capacitance matrix, fast Poisson solvers, conjugate gradients

CR Categories: 5.17
Language: FORTRAN

DESCRIPTION

This algorithm provides an approximate solution to the Helmholtz equation

$$-\Delta u + cu = g \quad \text{in } \Omega$$

with a Dirichlet boundary condition

$$u = f \quad \text{on } \Gamma, \text{ the boundary of } \Omega.$$

Here Ω , a three-dimensional bounded region, c , an arbitrary real constant (positive, negative, or zero), and the functions f and g are specified by the user. The Laplace operator Δ is in Cartesian coordinates.

A second-order accurate finite-difference method is used to discretize the Helmholtz equation. The resulting linear system of equations is reduced to a capacitance matrix equation that is solved approximately by a conjugate gradient method. We sketch the basic ideas below, but a detailed discussion of this and similar methods can be found in [1].

To perform the discretization, the region Ω is embedded in a cube and a uniform rectangular finite-difference grid is imposed. A simple seven-point difference approximation is used for all mesh points except those that are in Ω and are near the boundary Γ . For these boundary neighbors, second-order accurate equations incorporating the boundary data are used. The resulting difference

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scheme is known as the Shortley-Weller method. Thus the discrete system of equations has a matrix that differs from that for a Helmholtz problem on the cube only in those rows that correspond to points near Γ . We take advantage of this by reformulating the problem as one of dimension equal to the number of boundary neighbors rather than the number of mesh points in the region. The resulting linear system is the capacitance matrix equation. In our implementation this reduced equation is then solved using an iterative algorithm, the conjugate gradient method. A special scaling method, based on so-called discrete dipoles, is used to enhance the convergence. A fast Poisson solver on the cube is one of the components necessary to evaluate the product of the capacitance matrix with a given vector.

Let NX , NY , and NZ be the numbers of mesh points in the cube along the three coordinate axes, where NX and NY are powers of 2. We denote the number of mesh points in Ω with at least one of their six nearest neighbors on or outside the boundary Γ by $IPP1 + IPP2$, where $IPP2$ points have two such neighbors along at least one coordinate mesh line. The program requires as input certain scalar parameters, the coordinates of each of the $IPP1 + IPP2$ points and their distances to the boundary along mesh lines, the values of the Dirichlet data f at points of intersection of mesh lines and the boundary Γ , and the value of the function g at mesh points in Ω . The user communicates with the package through the subroutine HELM3D. A complete description of the input parameters is given in the comments at the beginning of this subroutine.

HELM3D controls the conjugate gradient iteration and calls upon a fast solver (CUBE) and subroutines UTAMLT, UTATRN, BNDRY, VMULT, and VTRANS to perform the necessary matrix-vector products. CUBE solves the Helmholtz equation on a cube using fast discrete Fourier transform routines RFORT and FORT to reduce the systems to tridiagonal form. The resulting linear systems are solved by a Toeplitz method, and then an inverse Fourier transform is performed. RFORT and FORT were provided by Dr. W. Proskowski, who has modified a code written by Dr. J. Cooley. HELM3D also employs an error-checking module HELMCK to check the input data. It diagnoses errors in the integer parameters, missing boundary points, and inconsistencies in the given boundary data.

The program requires two arrays of dimension $NX \times NY \times NZ$ (one if $g = 0$), four integer and six real arrays of dimension $IPP1 + 2 * IPP2$, and one real array of dimension $\max(IPP1 + 2 * IPP2, NX * NZ, NY * NZ)$. Each conjugate gradient iteration requires time proportional to $NX * NY * NZ * \text{LOG}(NX * NY)$, and the number of iterations will usually be small unless a value of c is used that makes the discrete Helmholtz operator almost singular. Double precision is required on machines with a short word length.

Data on timing for many sample problems are given in [1]. As an example, a discrete Laplace problem on a cube with a sphere cut out of it having 10464 mesh points and embedded in a cube of dimension $32 \times 32 \times 24$ required 84K words of storage and 171 seconds on a CDC 6600 (FTN compiler, $OPT = 2$) to find a solution of the linear system with a maximum error equal to 0.55×10^{-3} of the maximum value of the solution.

A sample driver is included with the algorithm. Possible enhancements to the algorithm are discussed in [1].

REFERENCES

1. O'LEARY, D. P., AND WIDLUND, O. Capacitance matrix methods for the Helmholtz equation on general three dimensional regions. Courant Inst. Rep. COO-3077-155, New York, Oct. 1978; also *Math. Comp.* 33 (1979), 849-879.

ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```

SUBROUTINE HELM3D (MODE, W, GG, NXDIM, NYDIM, NZDIM, IPP1, IPP2, DELTA, NNX
1, NNY, NNZ, NIPDIM, NAPDIM, ICCORD, INDORD, CC, NIT, EPS, S, R, P, AP, IER)
INTEGER MODE, NXDIM, NYDIM, NZDIM, IPP1, IPP2, NNX, NNY, NNZ, NIPDIM, ICCORD
1(3, NIPDIM), INDORD(NIPDIM), NIT, IER
REAL W(NXDIM, NYDIM, NZDIM), GG(NXDIM, NYDIM, NZDIM), DELTA(3, NIPDIM), CC
1, EPS, S(NIPDIM), R(NIPDIM), P(NIPDIM), AP(NAPDIM)

```

```

C
C   THIS PROGRAM WAS DEVELOPED BY DIANNE P O'LEARY AND OLOF WIDLUND.
C   THIS IS VERSION. MD2,  OCTOBER, 1979.
C

```

```

C   THIS PROGRAM SOLVES THE DIRICHLET PROBLEM FOR THE
C   HELMHOLTZ EQUATION OVER A GENERAL BOUNDED 3 DIMENSIONAL
C   REGION IMBEDDED IN A UNIT CUBE
C

```

```

C   -W   - W   - W   + CC*W = G1   IN THE REGION
C   XX   YY   ZZ
C

```

```

C   W = F                               ON THE BOUNDARY
C

```

```

C   WHERE F AND G1 ARE GIVEN FUNCTIONS OF X, Y, AND Z, AND CC IS
C   A REAL CONSTANT. THE BOUNDARY IS ARBITRARY. THE PROGRAM
C   PROVIDES A SOLUTION OF THE WELL KNOWN SHORTLEY-WELLER
C   APPROXIMATION OF THE DIFFERENTIAL EQUATION. THE MESH IS UNIFORM
C   IN EACH COORDINATE DIRECTION AND A SIMPLE SEVEN POINT FORMULA
C   IS USED FOR INTERIOR MESH POINTS. A CAPACITANCE MATRIX
C   METHOD, WITH DISCRETE DIPOLES, IS USED. THE CAPACITANCE
C   MATRIX EQUATION IS FORMULATED AS A LEAST SQUARES PROBLEM
C   AND SOLVED USING THE CONJUGATE GRADIENT METHOD.
C

```

REFERENCES:

```

C   O'LEARY AND WIDLUND, CAPACITANCE MATRIX METHODS
C   FOR THE HELMHOLTZ EQUATION ON GENERAL 3-DIMENSIONAL
C   REGIONS, NYU-DOE REPORT COO-3077-155, OCTOBER, 1978;
C   MATH. COMP. 33, 1979 849-880.
C

```

```

C   PROSKUROWSKI AND WIDLUND, MATH. COMP. 30, 1976 443-469.
C   ALSO NYU-DOE REPORT.
C

```

```

C   PROSKUROWSKI, LAWRENCE BERKELEY LAB REPORTS AND
C   "NUMERICAL SOLUTION OF HELMHOLTZ'S EQUATION BY
C   IMPLICIT CAPACITANCE MATRIX METHODS," ACM TRANS.
C   ON MATH. SOFTWARE 5, 1979 36-49.
C

```

```

C   SHIEH, MRC-WISCONSIN REPORTS AND NUMER. MATH. 29
C   1978 307-327.
C

```

MACHINE DEPENDENT FEATURES:

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C   THIS PROGRAM SHOULD BE CONVERTED TO DOUBLE PRECISION
C   IF IT IS TO BE USED ON COMPUTERS WITH SHORT WORD
C   LENGTH ,SUCH AS IBM 360/370.
C

```

GENERAL DESCRIPTION OF THE PARAMETERS:

```

C   INTEGER VALUES:
C   DIMENSIONS OF ARRAYS (NXDIM, NYDIM, NZDIM,
C   NIPDIM, NAPDIM)
C   NUMBER OF MESH POINTS IN CUBE (NNX X NNY X
C   NNZ)
C   NUMBER OF POINTS IN REGION ADJACENT TO BOUNDARY
C   (IPP1, IPP2)
C   MAXIMUM NUMBER OF ITERATIONS ALLOWED (NIT)
C   ERROR CODE (IER)
C   CODE TO CONTROL PROGRAM OPTIONS (MODE)
C

```

```

C   REAL VALUES:
C   HELMHOLTZ CONSTANT (CC)
C   CONVERGENCE TOLERANCE (EPS)
C

```

```

C
C   INTEGER ARRAYS:
C       COORDINATES OF POINTS IN REGION ADJACENT TO
C       BOUNDARY ('IRREGULAR POINTS') (ICoord)
C       WORK SPACE (INDORD)
C
C   REAL ARRAYS:
C       G1 VALUES (GG)
C       BOUNDARY VALUES (R, P, AP)
C       DISTANCES FROM IRREGULAR POINTS TO BOUNDARY
C       (DELTA)
C       WORK SPACE (W, S)
C
C   TOTAL ARRAY SPACE NEEDED:
C       REAL:  2  NXDIM * NYDIM * NZDIM (1 IF G1=0)
C              5  NIPDIM
C              1  NAPDIM
C       INTEGER:
C              4  NIPDIM
C
C       WHERE  NXDIM .GE. NNX,  NYDIM .GE. NNY,
C              NZDIM .GE. NNZ,
C              NIPDIM .GE. IPP1 + 2 * IPP2,
C              NAPDIM .GE. MAX (IPP1+2*IPP2, NNX*NNZ,
C                              NNY*NNZ)
C
C   NOTE:
C
C       IN THIS DOCUMENTATION, NN REFERS TO NNX, NNY, OR NNZ
C       AS APPROPRIATE, AND SIMILARLY H REFERS TO HX, HY, OR HZ.
C       THE MESH POINT (X,Y,Z) IS SAID TO HAVE 6 NEIGHBORS:
C       (X+HX,Y,Z), (X-HX,Y,Z), (X,Y+HY,Z), (X,Y-HY,Z),
C       (X,Y,Z+HZ), AND (X,Y,Z-HZ).
C       A MESH POINT IS CALLED IRREGULAR IF IT IS IN THE INTERIOR OF
C       THE REGION AND AT LEAST ONE OF ITS SIX NEIGHBORS IS ON OR
C       OUTSIDE THE BOUNDARY.
C
C   ON INPUT . . .
C   -- MODE =  1 IF THE REGION HAS BEEN CHANGED FROM THE PREVIOUS CALL
C              AND G1=0
C              2 IF THE REGION HAS BEEN CHANGED FROM THE PREVIOUS CALL
C              AND G1 IS NONZERO
C              3 IF THE REGION IS THE SAME AS ON THE PREVIOUS CALL
C              AND G1=0
C              4 IF THE REGION IS THE SAME AS ON THE PREVIOUS CALL
C              AND G1 IS NONZERO
C              5 IF THE PROBLEM IS THE SAME AS ON THE PREVIOUS CALL,
C              G1=0, AND THE ONLY CHANGE IS THAT EPS AND/OR NIT
C              MAY HAVE BEEN CHANGED
C              6 IF THE PROBLEM IS THE SAME AS ON THE PREVIOUS CALL,
C              G1 IS NONZERO, AND THE ONLY CHANGE IS THAT EPS
C              AND/OR NIT MAY HAVE BEEN CHANGED
C       IF MODE = 3,4,5, OR 5 DELTA, ICOORD, INDORD, NXDIM,
C       NYDIM, NZDIM, NNX, NNY, NNZ, IPP1, AND IPP2 MUST BE
C       UNCHANGED FROM THE PREVIOUS CALL. THE CURRENT VALUE OF S
C       WILL BE USED AS THE INITIAL GUESS FOR THE DIPOLE STRENGTHS.
C       (S=0 WILL BE USED IF MODE=1 OR 2.)
C       TO IMPROVE THE ACCURACY OF A PREVIOUSLY CALCULATED SOLUTION,
C       USE MODE=5 OR MODE=6 IF ROUND OFF IS NOT SUSPECTED. IF
C       ROUND OFF IS SUSPECTED, REINITIALIZE THE BOUNDARY VALUES IN R,
C       AP, AND P, AND USE MODE = 3 TO FORCE THE RESIDUAL TO BE
C       RECOMPUTED; IF G1 IS NONZERO, ADD GG TO THE SOLUTION
C       RETURNED BY THE SUBROUTINE.
C
C   -- W(NXDIM, NYDIM, NZDIM) IS UNINITIALIZED.
C   -- GG(NXDIM, NYDIM, NZDIM) INITIALIZED TO G1*HZ*HZ IN THE
C       REGION, WITH ARBITRARY VALUES OUTSIDE.
C       FOR I=1,...,NNX, J=1,...,NNY, AND K=1,...,NNZ,
C       GG(I,J,K) CORRESPONDS TO G1*((I-1)*HX, (J-1)*HY, (K-1)*HZ)*HZ**2.
C       IF MODE = 1, 3 OR 5, G1 MAY BE A DUMMY ARRAY (I.E.,
C       IT NEED NOT BE DIMENSIONED BY THE CALLING PROGRAM).
C
C   -- IPP1 IS THE NUMBER OF IRREGULAR POINTS WITH AT LEAST 1
C       INTERIOR NEIGHBOR IN EACH DIRECTION X, Y, AND Z.

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```

C      -- IPP2 IS THE NUMBER OF IRREGULAR POINTS WHICH, ALONG
C      AT LEAST ONE DIRECTION, HAVE TWO EXTERIOR NEIGHBORS.
C
C      IN THE EXCEPTIONAL CASE WHEN IPP1+IPP2 EQ. 0, THE ROUTINE
C      WILL SOLVE THE PROBLEM ON THE WHOLE CUBE WITH THE
C      BOUNDARY CONDITIONS#
C      G1(X,Y,Z) = 0      Z .LT. 0      OR Z .GT. 1
C      W(0,Y,Z) = W(1,Y,Z) AND W(X,0,Z) = W(X,1,Z)
C      W(X,Y,0)=0 AND W(X,Y,Z) BOUNDED FOR ALL Z.
C      ARRAY GG MUST BE INITIALIZED TO G1*HZ*HZ AND MODE = 2.
C      W MAY BE A DUMMY ARRAY. THE ANSWER WILL BE STORED
C      IN THE ARRAY GG IN THIS CASE.
C
C      -- DELTA(3, NIPDIM) RECORDS + OR - DISTANCE TO BOUNDARY
C      FROM IRREGULAR POINT L IN THE X, Y, AND Z
C      DIRECTIONS (3*IPP1 + 6*IPP2 VALUES). THESE DISTANCES
C      ARE EXPRESSED AS MULTIPLES OF THE MESH SPACING: I.E.,
C      IF A DELTA HAS THE VALUE Q, THE DISTANCE IS Q*H.
C      THERE ARE THREE DELTAS FOR EACH OF THE IPP1 POINTS
C      FOR L=1, IPP1,
C      DELTA(1,L) = SHORTER DISTANCE TO BOUNDARY ALONG X DIRECTION
C      DELTA(2,L) = SHORTER DISTANCE TO BOUNDARY ALONG Y DIRECTION
C      DELTA(3,L) = SHORTER DISTANCE TO BOUNDARY ALONG Z DIRECTION
C      THERE ARE SIX DELTAS FOR EACH OF THE IPP2 POINTS,
C      FOR L=1, IPP2 , LL=IPP1+2*L-1,
C      DELTA(1,LL) AND DELTA(1,LL+1) ARE THE DISTANCES TO THE
C      BOUNDARY ALONG THE POSITIVE AND NEGATIVE X DIRECTIONS
C      DELTA(2,LL) AND DELTA(2,LL+1) ARE THE DISTANCES TO THE
C      BOUNDARY ALONG THE POSITIVE AND NEGATIVE Y DIRECTIONS
C      DELTA(3,LL) AND DELTA(3,LL+1) ARE THE DISTANCES TO THE
C      BOUNDARY ALONG THE POSITIVE AND NEGATIVE Z DIRECTIONS
C      THE PROGRAM WILL INTERCHANGE DELTAS IF NECESSARY SO THAT
C      FOR L=1, IPP2 , LL=IPP1+2*L-1,
C      ABS(DELTA(S,LL)) LE. ABS(DELTA(S,LL+1)).
C      NO DELTA CAN BE SO CLOSE TO 0 AS TO CAUSE OVERFLOW
C      UPON DIVISION BY A PRODUCT OF TWO DELTAS. SUCH SMALL
C      DELTAS SHOULD BE AVOIDED BY CHANGING THE REGION
C      SLIGHTLY OR BY SHIFTING IT INSIDE THE CUBE OR BY
C      USING ANOTHER MESH SIZE.
C
C      -- NNX, NNY, NNZ ARE THE NUMBER OF MESH POINTS IN THE X, Y, AND Z
C      DIRECTIONS.
C      MAX(NNX,NNY) MUST BE .LE. 256 UNLESS THE ERROR CHECK IN
C      HELMCK AND THE DIMENSIONS OF IB AND S IN COMMON FFT
C      (SUBROUTINES CUBE,RFORT AND FORT) ARE CHANGED.
C      THE MESH SPACINGS WILL BE CALCULATED TO BE
C      HX = 1 / NNX
C      HY = 1 / NNY
C      HZ = 1 / (NNZ - 1)
C      NNX AND NNY MUST BE POWERS OF 2 AND .GE. 8 UNLESS
C      THE FFT ROUTINES RFORT AND FORT ARE REPLACED.
C
C      -- NIPDIM, THE DIMENSION OF THE ONE DIMENSIONAL ARRAYS,
C      MUST BE .GE. IPP1+2*IPP2.
C      -- NAPDIM , THE DIMENSION OF AP , MUST
C      BE .GE. MAX(IPP1+2*IPP2, NNX*NNZ, NNY*NNZ ).
C      -- ICOORD(3,NIPDIM) RECORDS THE 3*(IPP1+IPP2) INDICES OF
C      THE IRREGULAR POINTS. THESE INDICES MUST LIE BETWEEN
C      2 AND NN-1 INCLUSIVE.
C      FOR L = 1, IPP1
C      THE L-TH COLUMN OF ICOORD
C      GIVES THE INDICES CORRESPONDING
C      TO DATA IN THE L-TH COLUMNS OF
C      DELTA, P, R, AND AP.
C      FOR L = 1, IPP2 , LL = IPP1 + 2 * L - 1
C      THE (IPP1+L)-TH COLUMN OF ICOORD
C      GIVES THE INDICES CORRESPONDING TO
C      DATA IN THE LL-TH AND (LL+1)-TH
C      COLUMNS OF DELTA, P, R, AND AP.
C      -- INDORD (NIPDIM) IS UNINITIALIZED. THE PROGRAM WILL
C      RECORD A CODE (1-6) FOR THE ORDER OF THE DELTAS.
C      -- CC IS THE CONSTANT IN THE HELMHOLTZ EQUATION.
C      -- NIT IS THE MAXIMUM NUMBER OF CONJUGATE GRADIENT ITERATIONS
C      ALLOWED.
C      -- EPS IS THE TOLERANCE FOR THE EUCLIDEAN NORM OF
C      THE CAPACITANCE EQUATION RESIDUAL DIVIDED BY THE

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```

C          S Q R T OF THE DIMENSION OF THIS VECTOR.
C          T T T
C          RESIDUAL = C U F - C C S WHERE C = U AGV .
C          IT IS DIFFICULT TO GIVE A RELIABLE RULE OF
C          THUMB FOR THE CHOICE OF EPS. FOR MANY PROBLEMS
C          ONE TENTH OF THE DESIRED ACCURACY FOR THE
C          SOLUTION OF THE ORIGINAL DISCRETE PROBLEM IS A
C          SUITABLE VALUE. A SMALLER TOLERANCE IS REQUIRED
C          WHEN THE DISCRETE HELMHOLTZ OPERATOR IS CLOSE
C          TO SINGULAR.
C          -- S , P , R ARE OF DIMENSION NIPDIM .
C          AP IS OF DIMENSION NAPDIM
C          S IS UNINITIALIZED IF MODE = 1 OR 2.
C          IF MODE .LT. 5 , FOR L=1, IPP1+2*IPP2,
C          R(L) = F(X+DELTA(1,L)*HX, Y, Z)
C          P(L) = F(X, Y+DELTA(2,L)*HY, Z)
C          AP(L) = F(X, Y, Z+DELTA(3,L)*HZ)
C          WHERE X, Y, AND Z ARE THE COORDINATES OF THE
C          IRREGULAR POINT CORRESPONDING TO THE DELTAS.
C          THE VALUES OF R , P , AND AP ARE NOT USED IN THE
C          COMPUTATION IF THE ABSOLUTE VALUE OF THE CORRESPONDING
C          DELTA IS GREATER THAN 1.
C
C          --- IER IS UNINITIALIZED. THE PROGRAM WILL RECORD AN ERROR
C          CODE (0-3).
C          THE USE OF DISCRETE DIPOLES IMPOSES A MILD RESTRICTION
C          ON THE GEOMETRY OF THE REGION. THE THREE MESH POINTS, OBTAINED
C          BY STEPPING FROM AN IRREGULAR POINT IN THE DIRECTION OF THE
C          SMALLEST MAGNITUDE DELTA, FROM THIS NEW POINT IN
C          THE DIRECTION OF THE MEDIUM, AND FROM THERE IN THE DIRECTION
C          OF THE LARGEST MUST NOT BE INTERIOR POINTS OF THE REGION.
C          ASSOCIATED WITH AN IRREGULAR POINT WHICH HAS AT
C          LEAST 2 EXTERIOR NEIGHBORS IN SOME MESH
C          DIRECTION ARE TWO COLUMNS OF THE ARRAY DELTA
C          THE DELTA'S RELEVANT TO THIS TEST ARE THE
C          SMALLER IN MAGNITUDE OF THE TWO POSSIBLE
C          CHOICES IN EACH COORDINATE DIRECTION. IF THE
C          RESTRICTION IS VIOLATED, A SUBROUTINE HELMCK WILL RETURN AN
C          ERROR FLAG IER = 2. A REFINEMENT OF THE MESH OR
C          A SLIGHT SHIFT OF THE REGION IN THE UNIT CUBE MIGHT
C          RESOLVE THE PROBLEM.
C
C          ON OUTPUT . . .
C          W WILL CONTAIN VALUES OF THE SOLUTION INSIDE THE
C          REGION AND USELESS VALUES OUTSIDE AND ON THE
C          BOUNDARY.
C          S WILL RECORD DIPOLE STRENGTHS. THIS IS THE SOLUTION
C          VECTOR OF THE CAPACITANCE MATRIX EQUATION.
C          R WILL BE THE RESIDUAL OF THE CAPACITANCE EQUATION.
C
C          P , AP , AND GG WILL BE CHANGED, AND THE DELTAS MAY
C          BE REORDERED AS INDICATED ABOVE.
C
C          ERROR RETURNS:
C          IER=0 NO ERROR
C          =1 ERROR IN INTEGER PARAMETER
C          =2 ERROR IN ICOORD OR VIOLATION OF DIPOLE
C          RESTRICTION OR IRREGULAR POINT MISSING
C          =3 TOO MANY CONJUGATE GRADIENT ITERATIONS
C          WITHOUT CONVERGENCE. ANSWER DOES NOT
C          HAVE THE REQUESTED ACCURACY.
C
C          AFTER EACH ITERATION, THE FOLLOWING INFORMATION IS PRINTED:
C          -- THE CONJUGATE GRADIENT PARAMETERS ALPHA AND BETA.
C          THIS INFORMATION COULD BE USED TO ESTIMATE THE
C          CONDITION NUMBER OF THE CAPACITANCE MATRIX.
C          -- THE EUCLIDEAN NORM OF THE RESIDUAL OF THE
C          CAPACITANCE MATRIX EQUATION.
C          T T T
C          THE RESIDUAL=C U F-C C S WHERE C= U AGV.
C
C          THE ROLES OF THE SUBROUTINES:
C          HELM3D CONTROLS THE CONJUGATE GRADIENT ITERATION.
C          HELMCK CHECKS THE INPUT DATA FOR CORRECTNESS.
C          VMULT USES THE DIPOLE STRENGTHS IN A NIPDIM ARRAY TO

```

```
C          SET UP THE DIPOLES IN A 3 DIMENSIONAL ARRAY.
C          THIS SUBROUTINE THUS DEFINES A LINEAR MAPPING
C          FROM A SPACE OF 1-DIMENSIONAL ARRAYS TO A SPACE
C          OF 3-DIMENSIONAL ARRAYS.
C          VTRANS DEFINES THE TRANSPOSE OF THE MAPPING DEFINED
C          BY VMULT.
C          UTAMLT MAPS 3-DIMENSIONAL ARRAYS INTO 1-DIMENSIONAL
C          ARRAYS BY USING A FINITE DIFFERENCE FORMULA WHICH
C          CORRESPONDS TO A PART OF THE SHORTLEY-WELLER
C          APPROXIMATION. THE REMAINING PART IS HANDLED BY
C          BNDRY.
C          UTATRN DEFINES THE TRANSPOSE OF THE MAPPING DEFINED BY
C          UTAMLT.
C          BNDRY PROCESSES THE DIRICHLET DATA AND THE VALUES OF G1
C          CLOSE TO THE BOUNDARY, PRODUCING U(TRANSPOSE)F FOR
C          USE IN THE RIGHT HAND SIDE OF THE CAPACITANCE EQUATION
C          CUBE SOLVES THE HELMHOLTZ EQUATION OVER A CUBE USING A
C          FOURIER-TOEPLITZ ALGORITHM.
C          RFORT IS A FAST FOURIER TRANSFORM ROUTINE DUE TO
C          W. PROSKUROWSKI WHO REVISED A CODE WRITTEN BY J. COOLEY.
C          IT IS USED BY SUBROUTINE CUBE.
C          FORT IS A SUBROUTINE CALLED BY RFORT.
```


ALGORITHM 573

NL2SOL—An Adaptive Nonlinear Least-Squares Algorithm [E4]

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and

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Key Words and Phrases: unconstrained optimization, nonlinear least squares, nonlinear regression, quasi-Newton methods, secant methods

CR Categories: 5.14, 5.5

Language: FORTRAN

1. PURPOSE

Given a continuously differentiable function (residual vector) $R(x) = (R_1(x), R_2(x), \dots, R_n(x))^T$ of p parameters $x = (x_1, x_2, \dots, x_p)^T$, NL2SOL attempts to find a parameter vector x^* that minimizes the sum-of-squares function $F(x) = \sum_{i=1}^n R_i(x)^2$.

2. METHOD

Reference [1] explains the algorithm realized by NL2SOL in detail. The algorithm amounts to a variation on Newton's method in which part of the Hessian matrix is computed exactly and part is approximated by a secant (quasi-Newton) updating method. Once the iterates come sufficiently close to a local solution, they usually converge quite rapidly. To promote convergence from poor starting guesses, NL2SOL uses a model/trust-region technique along with an adaptive choice of the model Hessian. Consequently, the algorithm sometimes reduces to a Gauss-Newton or Levenberg-Marquardt method. On large residual problems (in which $F(x^*)$ is large), however, NL2SOL often works much better than these methods.

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3. DESCRIPTION

3.1 Calling Sequence

CALL NL2SOL (N, P, X, CALCR, CALCJ, IV, V, UIPARM, URPARM, UFPARM)

Note: NL2SOL is written in American National Standard FORTRAN (1966) and the comments below assume that the calling program is also written in FORTRAN. These comments refer to the single-precision version of NL2SOL. In the double-precision version, all quantities termed REAL below are actually DOUBLE PRECISION.

- N** (input INTEGER) is the number of components in the residual vector R .
- P** (input INTEGER) is the number of parameters on which R depends.
- X** (I/O REAL array of length P) on input is an initial guess at the desired solution x^* . When NL2SOL returns, X contains the best parameter estimate found so far.
- CALCR** (input subroutine) computes the residual vector $R = R(X)$ when invoked by
 CALL CALCR(N, P, X, NF, R, UIPARM, URPARM, UFPARM)
 When CALCR is called, NF is the invocation count for CALCR; it is included for possible use with CALCJ. If X is out of bounds (e.g., if $R(X)$ would overflow), then CALCR should set NF to 0, which will cause a shorter step to be attempted. CALCR should not change N, P, or X and should be declared EXTERNAL in the calling program. R should be declared REAL R(N).
- CALCJ** (input subroutine) computes the Jacobian matrix $J = J(X)$ of first partials, $J_{ij} = \partial J_i(X) / \partial x_j$, when invoked by
 CALL CALCJ(N, P, X, NF, J, UIPARM, URPARM, UFPARM)
 When CALCJ is called, NF is the invocation count for CALCR at the time when $R(X)$ was evaluated. The X passed to CALCJ is usually the one passed to CALCR on either its most recent invocation or the one prior to it. Thus if CALCR saves intermediate results for use by CALCJ, then it is possible to tell from NF whether they are valid for the current X (or which copy is valid if two are kept). If J cannot be computed at X, then CALCJ should set NF to 0. CALCJ should not change N, P, or X and should be declared EXTERNAL in the calling program. J should be declared REAL J(N, P).
- IV** (I/O INTEGER array of length P + 60) on input contains certain values (such as limits on the number of iterations and calls on CALCR) that control the behavior of NL2SOL and on output contains various counts and other items of interest: see Sections 3.3 and 3.4. If IV(1) = 0 on input, then default values are supplied for the input components of both IV and V. The caller may supply nondefault values for selected components of IV and V by calling DFAULT(IV, V) and then assigning the appropriate nondefault values before calling NL2SOL.
- V** (I/O REAL array of length $93 + N(P + 3) + P(3P + 33)/2$) on input contains certain values (such as convergence tolerances) that control the behavior of NL2SOL and on output contains various items of interest (such as $F(X)$ and $R(X)$): see Sections 3.5 and 3.15.
- UIPARM** (INTEGER array of length determined by the caller) is passed without change to CALCR and CALCJ and may be used by them in any way that the caller may find convenient.
- URPARM** (REAL array of length determined by the caller), like UIPARM, is passed without change to CALCR and CALCJ.

UFPARM (subroutine), like UIPARM, is passed without change (and without having been invoked) to CALCR and CALCJ. If there is no need for such a subroutine, then on many systems it suffices to pass an arbitrary variable or constant for UFPARM. But if an actual subroutine is passed, then it must be declared EXTERNAL in the calling program.

3.2 Example

Let $n = 3$, $p = 2$, and

$$R(x) = \begin{bmatrix} x_1^2 + x_2^2 + x_1x_2 \\ \sin x_1 \\ \cos x_2 \end{bmatrix}.$$

(This problem is due to Madsen [3].) The following FORTRAN code minimizes $F(x) = \frac{1}{2}R(x)^T R(x)$, starting from the initial guess $(3, 1)^T$, using a single-precision version of NL2SOL.

```

INTEGER IV(62), UI(1)
REAL V(147), X(2), UR(1)
EXTERNAL MADR, MADJ
X(1) = 3.0
X(2) = 1.0
IV(1) = 0
CALL NL2SOL (3, 2, X, MADR, MADJ, IV, V, UI, UR, MADR)
STOP
END
SUBROUTINE MADR (N, P, X, NF, R, UI, UR, UF)
INTEGER N, P, NF, UI(1)
REAL X(P), R(N), UR(1)
EXTERNAL UF
R(1) = X(1)**2 + X(2)**2 + X(1)*X(2)
R(2) = SIN(X(1))
R(3) = COS(X(2))
RETURN
END
SUBROUTINE MADJ (N, P, X, NF, J, UI, UR, UF)
INTEGER N, P, NF, UI(1)
REAL X(P), J(N, P), UR(1)
EXTERNAL UF
J(1, 1) = 2.0*X(1) + X(2)
J(1, 2) = 2.0*X(2) + X(1)
J(2, 1) = COS(X(1))
J(2, 2) = 0.0
J(3, 1) = 0.0
J(3, 2) = -SIN(X(2))
RETURN
END

```

The main program above passes MADR as CALCR and MADJ as CALCJ. No use is made of UIPARM, URPARM, or UFPARM in this simple example.

When the above is executed, NL2SOL prints the initial X vector, a summary of the iterations performed, the final X vector, and some statistics, including the final $F(X)$ and a covariance matrix. If REAL is changed to DOUBLE PRECISION and the above is run on an IBM 370 computer, then NL2SOL reports relative function convergence $(IV(1) = 4)$ —see Section 3.3) after 12 calls on MADR and MADJ and returns $X(1) = -0.155437$, $X(2) = 0.694564$, and $F(X) = 0.386600$.

If, say, we wanted to suppress the iteration summary, we could do so by replacing the statement $IV(1) = 0$ in the main program by

```

CALL DFAULT(IV, V)
IV(19) = 0

```

(See the description of IV(OUTLEV) in Section 3.4.)

3.3 Return Codes

When NL2SOL returns, IV(1) contains one of the following return codes:

- 3 = X-convergence. The scaled relative difference between the current parameter vector X and a locally optimal parameter x^* is very likely at most $V(XCTOL)$: see Section 3.5.
- 4 = relative function-convergence. The relative difference between the current function value and its locally optimal value is very likely at most $V(RFCTOL)$: see Section 3.5.
- 5 = both X and relative function-convergence, that is, the conditions for IV(1) = 3 and IV(1) = 4 both hold.
- 6 = absolute function-convergence. The current function value (half the sum of squares) is at most $V(AFCTOL)$: see Section 3.5.
- 7 = singular convergence. The Hessian near the current X appears to be singular or nearly so, and a step of scaled length at most $V(LMAX0)$ is unlikely to yield a relative function decrease of more than $V(RFCTOL)$. This means that the model is over-specified (i.e., contains too many parameters), at least near X . It is possible that a different starting guess would lead NL2SOL to find an X giving a smaller $F(X)$ and strong convergence (IV(1) = 3, 4, 5, or 6).
- 8 = false convergence. The iterates appear to be converging to a noncritical point. This may mean that the false convergence tolerance ($V(XFTOL)$)—see Section 3.5) is too large, that the convergence tolerances ($V(AFCTOL)$, $V(RFCTOL)$, $V(XCTOL)$) are too small for the accuracy to which CALCR and CALCJ compute R and J , that there is an error in computing the Jacobian matrix J , or that R is discontinuous near X .
 If the NPRELDF value printed in the summary statistics (or in the iteration summary for the final iteration) is negative and not too much larger than $V(RFCTOL)$ in absolute value, then $V(RFCTOL)$ is too small and singular convergence would be detected if $V(RFCTOL)$ were increased above $|NPRELDF|$: see Sections 3.5 and 3.11.
- 9 = function evaluation limit reached without other convergences: see IV(MXFCAL) in Section 3.4.
- 10 = iteration limit reached without other convergence: see IV(MXITER) in Section 3.4.
- 11 = STOPX returned .TRUE. (external interrupt): see Section 3.14.
- 13 = $F(X)$ cannot be computed at the initial X .
- 14 = bad parameters passed to ASSESS. (This should not occur.)
- 15 = the Jacobian could not be computed at X (see CALCJ above).
- 16 = N or P (or parameter NN to NL2ITR) out of range: $P < 0$ or $N < P$ or $NN < N$.
- 17 = a restart was attempted with N , P , or parameter NN to NL2ITR changed: see Section 3.7.
- 18 = IV(INITS) out of range: see Section 3.4.
- 19-45 = $V(IV(1))$ is out of range.
- 50 = IV(1) was out of range when NL2SOL (or NL2SNO or NL2ITR) was called.
- 87 ... (86 + P) = JTOL(IV(1)-86), that is, $V(IV(1))$, is not positive: see $V(DFAC)$ in Section 3.5.

Just before NL2SOL returns, a brief description of the return code is printed (unless all printing is turned off by IV(PRUNIT) = 0).

3.4 IV Values

3.4.1 IV Input Values (Supplied by DFAULT)

IV(1) . . . IV(1) should have a value between 0 and 12 when NL2SOL is called. 0 and 12 both mean that this is a fresh start; 0 means DFAULT(IV, V) should be invoked to supply default values to the input components of IV and V, while 12 (the value that DFAULT assigns to IV(1)) means that the caller has already called DFAULT(IV, V) and has possibly changed some IV or V entries to nondefault values. IV(1) input values between 3 and 11 mean that NL2SOL should restart; see Section 3.7. Default = 12.

IV(COVPRT) . . . IV(14) = 1 means print a covariance matrix at the solution. This matrix is computed as IV(COVREQ) dictates just before a return with IV(1) = 3, 4, 5, or 6. IV(COVPRT) = 0 means skip this printing. Default = 1.

IV(COVREQ): IV(15) \neq 0 means compute a covariance matrix before a return with IV(1) = 3, 4, 5, or 6. In this case, an approximate covariance matrix is obtained in one of several ways. Let $k = |IV(COVREQ)|$ and let $\sigma = 2F(X)/\max(1, N-P)$, where $2F(X)$ is the residual sum of squares. If $k = 1$ or 2, then a finite-difference Hessian approximation H is obtained. If H is positive-definite (or, for $k = 3$, if the Jacobian has full rank), then one of the following is computed:

$$k = 1 \Rightarrow \sigma H^{-1}(J^T J)H^{-1}$$

$$k = 2 \Rightarrow \sigma H^{-1}$$

$$k = 3 \Rightarrow \sigma (J^T J)^{-1}.$$

If IV(COVREQ) > 0, then both function and gradient values (calls on CALCR and CALCJ) are used in computing H (with step sizes determined by V(Delta0); see Section 3.5), while if IV(COVREQ) < 0, then only function values (calls on CALCR) are used (with step sizes determined by V(DLTFDC)). If IV(COVREQ) = 0, then no attempt is made to compute a covariance matrix (unless IV(COVPRT) = 1, in which case NL2SOL assumes IV(COVREQ) = 1 and NL2SNO assumes IV(COVREQ) = -1). See IV(COVMAT) below. Default = 1.

IV(DTYPE) . . . IV(16) tells how the scale vector D (see [1]) should be chosen. IV(DTYPE) > 0 means choose D as described below with V(DFAC). IV(DTYPE) \leq 0 means the caller has chosen D and has stored it in V starting at V(94 + 2N + P(3P + 31)/2). Default = 1.

IV(INITS) . . . IV(25) tells how the S matrix (see [1]) should be initialized. 0 means initialize S to all zeros and start with the Gauss-Newton model. 1 and 2 mean that the caller has stored the lower triangle of the initial S rowwise in V starting at V(87 + 2P). IV(INITS) = 1 means start with the Gauss-Newton model, while IV(INITS) = 2 means start with the augmented model; see [1]. Default = 0.

IV(MXFCAL) . . . IV(17) gives the maximum number of function evaluations (calls on CALCR, excluding those used to compute the covariance matrix and, in the case of NL2SNO, the Jacobian matrices) allowed. If this number does not suffice, then NL2SOL returns with IV(1) = 9. Default = 200.

IV(MXITER) . . . IV(18) gives the maximum number of iterations allowed. It also indirectly limits the number of gradient evaluations (calls on CALCJ) to IV(MXITER) + 1. If IV(MXITER) iterations do not suffice, then NL2SOL returns with IV(1) = 10. Default = 150.

IV(OUTLEV) . . . IV(19) controls the number and length of iteration summary lines printed (by ITSMRY). IV(OUTLEV) = 0 means do not print any summary lines. Otherwise print a summary line after each $|IV(OUTLEV)|$

iterations. Long summary lines are printed if $IV(OUTLEV) > 0$, short lines if $IV(OUTLEV) < 0$. See Section 3.11 for more details. Default = 1.

$IV(PARPRT) \dots IV(20) = 1$ means print any nondefault V values on a fresh start or any changed V values on a restart. $IV(PARPRT) = 0$ means skip this printing. Default = 1.

$IV(PRUNIT) \dots IV(21)$ is the output unit number on which all printing is done. $IV(PRUNIT) = 0$ means suppress all printing. (Setting $IV(PRUNIT)$ to 0 is the only way to suppress the one-line termination message printed before $NL2SOL$ returns.) Default = standard output unit (unit 6 on most systems); the default for $IV(PRUNIT)$ is actually $IMDCON(1)$; see Section 3.12.

$IV(SOLPRT) \dots IV(22) = 1$ means print the final X (the one returned), along with the final gradient and scale vector D . $IV(SOLPRT) = 0$ means skip this printing. Default = 1.

$IV(STATPR) \dots IV(23) = 1$ means print summary statistics upon returning. These consist of the function value (half the residual sum of squares) at X , the scaled relative size of the last step taken (see $V(RELDX)$ below), the number of function and gradient evaluations (calls on $CALCR$ and $CALCJ$, excluding any calls made only for computing covariance matrices), the relative function reductions predicted for the last step taken and for a Newton step (or perhaps a step of length bounded by $V(LMAX0)$ —see the descriptions of $PRELDF$ and $NPRELDF$ in Section 3.11 below), and, if an attempt was made to compute a covariance matrix, the number of calls on $CALCR$ and $CALCJ$ used in trying to compute the covariance matrix. $IV(STATPR) = 0$ means skip this printing. Default = 1.

$IV(X0PRT) \dots IV(24) = 1$ means print the initial X and scale vector D (on a fresh start only). $IV(X0PRT) = 0$ means skip this printing. Default = 1.

3.4.2 *IV Output Values of Primary Interest*

$IV(1) \dots IV(1)$ is the return code; see Section 3.3.

$IV(COVMAT) \dots IV(26)$ tells whether a covariance matrix was computed. If $IV(COVMAT) > 0$, then the lower triangle of the covariance matrix is stored row-wise in V , starting at $V(IV(COVMAT))$. If $IV(COVMAT) = 0$, then no attempt was made to compute a covariance matrix. If $IV(COVMAT) = -1$, then the finite-difference Hessian H was indefinite (or, for $|IV(COVREQ)| = 3$, the current Jacobian matrix is rank deficient): like singular convergence (see Section 3.3), this may mean that the model is overspecified (contains too many parameters). And if $IV(COVMAT) = -2$, then a successful finite-difference step could not be found for some component of X (i.e., $CALCR$ set NF to 0 for each of two trial steps).

Note that $IV(COVMAT)$ is reset to 0 after each successful step, so if such a step is taken after a restart, then the covariance matrix will be recomputed.

$IV(D) \dots IV(27)$ is the starting subscript in V of the current scale vector D .

$IV(G) \dots IV(28)$ is the starting subscript in V of the current least-squares gradient vector $J^T R$.

$IV(NFCALL) \dots IV(6)$ is the number of calls so far made on $CALCR$ (i.e., function evaluations, including those used in computing covariance matrices).

$IV(NFCOV) \dots IV(40)$ is the number of calls made on $CALCR$ when computing covariance matrices.

$IV(NGCALL) \dots IV(30)$ is the number of calls on $CALCJ$ (gradient evaluations) so far made, including those used in computing covariance matrices.

IV(NGCOV) . . . IV(41) is the number of calls made on CALCJ when computing covariance matrices.

IV(NITER) . . . IV(31) is the number of iterations performed.

IV(R) . . . IV(50) is the starting subscript in V of the residual vector R corresponding to the final X .

3.5 V Values of Primary Interest

Many of the V input components described here and in Section 3.15 must lie in certain intervals. If such a component lies outside the interval indicated for it below (or in Section 3.15) at the beginning of its description, then module PARCHK will print an error message (unless IV(PRUNIT) = 0) and will force NL2SOL to return immediately with IV(1) > 18.

Frequent reference is made below to two quantities: MACHEP and the scale vector D . MACHEP is the unit roundoff for the floating-point arithmetic being used—see Section 3.12. The scale vector D is the diagonal of the diagonal scale matrix D_k discussed in [1, Sections 5 and 7]; this scale matrix is denoted by $\text{diag}(D)$ below.

3.5.1 V Input Values of Primary Interest (Supplied by DFAULT)

V(AFCTOL) . . . V(31) > 0 is the absolute function convergence tolerance. If NL2SOL finds a point where the function value (half the sum of squares) is less than V(AFCTOL), and if NL2SOL does not return with IV(1) = 3, 4, or 5, then it returns with IV(1) = 6.

$$\text{Default} = \max\{10^{-20}, \text{MACHEP}^2\}.$$

V(DELTA0) . . . V(44) \in [MACHEP, 1] is a factor used in choosing the finite-difference step sizes used in computing covariance matrices when IV(COVREQ) = 1 or 2. For differences involving $X(i)$, step size

$$V(\text{DELTA0}) \cdot \max\{|X(i)|, 1/D(i)\} \cdot \text{sign}(X(i))$$

is used, where D is the current scale vector; see [1]. If this results in CALCR setting NF to 0, then -0.5 times this step is also tried. Default = $\text{MACHEP}^{1/2}$.

V(DFAC) . . . V(41) \in [0, 1] and the D0 and JTOL arrays (see V(D0INIT) and V(JTINIT)) are used in updating the scale vector D when IV(DTYPE) > 0. (D is initialized according to V(DINIT).) Let

$$D1(i) = \max\{[\text{JCNORM}(i)^2 + \max\{S_{ii}, 0\}]^{1/2}, V(\text{DFAC})D(i)\},$$

where JCNORM(i) is the 2-norm of the i th column of the current Jacobian matrix and S is the S matrix of [1]. If IV(DTYPE) = 1, then $D(i)$ is set to $D1(i)$ unless $D1(i) < \text{JTOL}(i)$, in which case $D(i)$ is set to $\max\{D0(i), \text{JTOL}(i)\}$. If IV(DTYPE) > 1, then D is updated during the first iteration as for IV(DTYPE) = 1 (after any initialization due to V(DINIT)) and is left unchanged thereafter. Default = 0.6.

V(DINIT) . . . V(38) ≥ -10 : if V(DINIT) ≥ 0 , then it is the value to which all components of the scale vector D are initialized during a fresh start. Default = 0.

V(DLTFDC) . . . V(40) \in [MACHEP, 1] helps choose the step sizes used in computing covariance matrices when IV(COVREQ) = -1 or -2 . For differences involving $X(i)$, the step size first tried is

$$V(\text{DLTFDC}) \cdot \max\{|X(i)|, 1/D(i)\},$$

where D is the current scale vector (see [1]). If this step is too big the first time it is tried, that is, if CALCR sets NF to 0, then -0.5 times this step is also tried. Default = $\text{MACHEP}^{1/3}$.

V(DLTFDJ) ... V(36) \in [MACHEP, 1] helps choose the step sizes used when NL2SNO computes a finite-difference approximation to the Jacobian matrix. For differences involving $X(i)$, the step size first tried is

$$V(DLTFDJ) \cdot \max\{|X(i)|, 1/D(i)\},$$

where D is the current scale vector (see [1]). If the first step is too big, that is, if CALCR sets NF to 0, then smaller steps are tried until the step size is shrunk below $1000 \cdot \text{MACHEP}$. Default = $\text{MACHEP}^{1/2}$.

V(D0INIT) ... V(37) ≥ 0 : if V(D0INIT) > 0 , it is the value to which all components of the D0 vector (see V(DFAC)) are initialized. If V(D0INIT) = 0, then it is assumed that the caller has stored D0 in V starting at V(P + 87). Default = 1.0.

V(JTINIT) ... V(39) ≥ 0 : if V(JTINIT) > 0 , it is the value to which all components of the JTOL array (see V(DFAC)) are initialized. If V(JTINIT) = 0, then it is assumed that the caller has stored JTOL in V starting at V(87). Default = 10^{-6} .

V(LMAX0) ... V(35) > 0 gives the maximum 2-norm allowed for $\text{diag}(D)$ times the very first step that NL2SOL attempts. It is also used in testing for singular convergence: if the function reduction predicted for a step of length bounded by V(LMAX0) is at most V(RFCTOL) $|F_0|$, where F_0 is the function value at the start of the current iteration, and if NL2SOL does not return with IV(1) = 3, 4, 5, or 6, then it returns with IV(1) = 7. Default = 100.

V(RFCTOL) ... V(32) \in [MACHEP, 0.1] is the relative function-convergence tolerance. If the current model predicts a maximum possible function reduction (see V(NREDUC)) of at most V(RFCTOL) $|F_0|$, where F_0 is the function value at the start of the current iteration, and if the last step attempted achieved no more than twice the predicted function decrease, then NL2SOL returns with IV(1) = 4 (or 5). Default = $\max\{10^{-10}, \text{MACHEP}^{2/3}\}$.

V(TUNER1) ... V(26) \in [0, 0.5] helps decide when to check for false convergence and to consider switching models. This is done if the actual function decrease from the current step is no more than V(TUNER1) times its predicted value. Default = 0.1.

V(XCTOL) ... V(33) \in [0, 1] is the X-convergence tolerance. If a Newton step (see V(NREDUC)) is tried that has V(RELDX) \leq V(XCTOL) and if this step yields at most twice the predicted function decrease, then NL2SOL returns with IV(1) = 3 (or 5). Default = $\text{MACHEP}^{1/2}$.

V(XFTOL) ... V(34) \in [0, 1] is the false-convergence tolerance. If a step is tried that gives no more than V(TUNER1) times the predicted function reduction and that has V(RELDX) \leq V(XFTOL), and if NL2SOL does not return with IV(1) = 3, 4, 5, 6, or 7, then it returns with IV(1) = 8. (See the description of V(RELDX) below.) Default = $100 \cdot \text{MACHEP}$.

V(*) ... DFAULT supplies to V a number of tuning constants, with which it should normally be unnecessary to tinker. See Section 3.15.

3.5.2 V Output Values of Primary Interest

V(DGNORM) ... V(1) = $\|\text{diag}(D)^{-1}g\|_2$, where g is the most recently computed gradient and D is the corresponding scale vector.

V(DSTNRM) ... V(2) = $\|\text{diag}(D)\Delta x\|_2$, where Δx is the most recently computed step and D is the current scale vector.

V(F) ... V(10) is the current function value (half the residual sum of squares).

V(F0) ... V(13) is the function value at the start of the current iteration.

V(NREDUC) . . . V(6), if positive, is the maximum function reduction possible according to the current model, that is, the function reduction predicted for a Newton step: $\Delta x = -H^{-1}g$, where $g = J^T R$ is the current gradient and H is the current Hessian approximation:

$$H = J^T J \quad \text{for the Gauss-Newton model}$$

$$H = J^T J + S \quad \text{for the augmented model.}$$

V(NREDUC) = 0 means H is not positive definite.

If V(NREDUC) < 0, then V(NREDUC) is used in the singular convergence test: It is the negative of the function reduction predicted for a step computed with a step bound of V(LMAX0).

V(PREDUC) . . . V(7) is the function reduction predicted (by the current quadratic model) for the current step. This (divided by V(F0)) is used in testing for relative function convergence.

V(RADIUS) . . . V(8) is the trust region radius (i.e., step bound) used for the last step tried.

V(RELDX) . . . V(17) is the scaled relative change in X caused by the current step, Δx , computed as

$$\max_i \{ |D(i)[X(i) - X_0(i)]| \} / \max_i \{ D(i)[|X(i)| + |X_0(i)|] \},$$

where $X = X_0 + \Delta x$.

3.6 Finite-Difference Jacobians: NL2SNO

Those who do not wish to code a subroutine CALCJ for (analytically) computing the Jacobian matrix may avoid doing so by calling NL2SNO instead of NL2SOL. NL2SNO computes an approximate Jacobian matrix by forward differences (using a step size determined by V(DLTFD)—see Section 3.5). The calling sequence for NL2SNO amounts to the one for NL2SOL with CALCJ omitted:

CALL NL2SNO (N, P, X, CALCR, IV, V, UIPARM, URPARM, UFPARM)

The parameters for NL2SNO are the same as the corresponding ones for NL2SOL, with the minor exception of IV(COVREQ): If IV(COVPR) = 1 and IV(COVREQ) = 0, then NL2SNO sets IV(COVREQ) to -1; otherwise, it sets IV(COVREQ) to -|IV(COVREQ)|. Thus NL2SNO uses function values only in computing covariance matrices and V(DELTA0) is not used.

3.7 Restarting

After any return with $3 \leq IV(1) \leq 11$, it is possible to change some of the IV and V input components (such as the convergence tolerances and the iteration and function evaluation limits) and call NL2SOL (or NL2SNO) again with IV(1) unchanged. This causes the algorithm to be resumed at the point where it was interrupted. (It is even possible to save IV, V, and X and then restart in a subsequent run.)

3.8 Scaling

Problems sometimes arise that are poorly scaled in the sense that the various components of X are expressed in widely differing units. With the default choice of the scale vector D (see V(DFAC) and the beginning of Section 3.5), the behavior of NL2SOL is largely insensitive to this kind of poor scaling. On well-scaled problems, the performance of NL2SOL can sometimes be improved by choosing D to be a vector of ones, that is, by setting IV(DTYPE) to 0 and V(DINIT) to 1.0. Occasionally it may also be worthwhile to fix $D(i)$, $1 \leq i \leq P$, at the 2-norm of the i th column of the initial Jacobian matrix by setting IV(DTYPE) to 2.

3.9 LMAX0: The Initial Step Bound

On some problems it is necessary to give $V(LMAX0) = V(35)$ a small value to prevent a disasterously large first step, one that might result in exponent overflow or arguments out of range to intrinsic functions. Even if no disaster occurs, if NL2SOL takes a number of function evaluations on the first iteration, then this number can be reduced on subsequent reruns by setting $V(LMAX0)$ to the value in the D*STEP column of the iteration summary for iteration 1.

3.10 Local Solutions

It can easily happen that NL2SOL only finds a local minimizer of the sum-of-squares function $F(X)$ and that a different starting guess would cause a point to be found at which F has a still smaller value. Except for cases where special conditions (such as convexity of the objective function) prevail, this shortcoming is shared by all minimization algorithm implementations.

3.11 Printed Output

Any printing is done by one of two modules: ITSMRY and PARCHK. PARCHK reports any V input components that are out of range and optionally lists any such components that have nondefault or changed values (on a fresh start or restart, respectively). ITSMRY does the remaining printing. Various IV input components control what printing is done; see Section 3.4.

If $IV(OUTLEV) > 0$, then ITSMRY produces an iteration summary which includes the following values: IT , the current iteration number; NF , the number of function evaluations (calls on CALCR), excluding any extra ones needed for computing covariance matrices and, in the case of NL2SNO, excluding the extra ones needed to compute finite-difference Jacobian matrices; F , the current function value (half the residual sum of squares); $RELDF$, the relative difference between the previous and the current function value (i.e., the difference in function values divided by the previous function value); $PRELDF$, the value of $RELDF$ predicted by the quadratic model used to compute the step just taken; $RELDX$, the relative change in X caused by the step just taken—see $V(RELDX)$ in Section 3.5; $MODEL$, a code that tells which models were used in choosing the current step (G = the Gauss-Newton model; S = the augmented model; $G-S$ means the Gauss-Newton model was tried first and a switch was then made to the augmented model; $S-G$, $G-S-G$, and $S-G-S$ have analogous meanings); $STPPAR$, the Marquardt parameter λ for the last step, Δx , computed: $\lambda > 0$ means Δx satisfies

$$[H + \lambda \text{diag}(D)^2]\Delta x = -g,$$

where H and g are the old Hessian approximation and gradient, respectively (the ones used in computing the step just taken); $SIZE$, the sizing factor used in updating the S matrix (see [1]); $D*STEP$, the 2-norm of $\text{diag}(D)$ times the step just taken (see $V(DSTNRM)$ in Section 3.5); and $NPRELDF$: if $NPRELDF > 0$, then it is the relative function reduction (i.e., value of $RELDF$) predicted for a full Newton step; if $NPRELDF = 0$, then the Hessian approximation failed to be positive definite; and if $NPRELDF < 0$, then it is the negative of the relative function reduction predicted for a step of length bounded by $V(LMAX0)$. These summary lines are produced every $IV(OUTLEV)$ iterations, and they are 118 characters long (including the carriage control character). If $IV(OUTLEV) < 0$, then short summary lines are produced every $-IV(OUTLEV)$ iterations; these lines are 79 characters long (55 if $IV(COVPR)$ = 0), and they include only the first six items listed above (i.e., IT , NF , F , $RELDF$, $PRELDF$, and $RELDX$).

3.12 Changing Computers

The NL2SOL distribution tape contains both single- and double-precision versions of the NL2SOL source code, so it should be unnecessary to change precisions. (On computers having only 32 or 36 bits per REAL word, double precision often gives better performance.)

Only the functions `IMDCON` and `RMDCON` contain machine-dependent constants. To change from one computer to another, it should suffice to change the `DATA` statements in these functions. The `DATA` statement in `IMDCON` sets `IMDCON(1)` to the output unit number that `DFAULT` supplies to `IV(PRUNIT)`. The machine-dependent `DATA` statement in `RMDCON` provides three values: `BIG`, `ETA`, and `MACHEP`. `BIG` is the largest floating-point number such that a FORTRAN program can compute $\text{SQRT}(0.999*\text{BIG})^{**2}$ (i.e., $\text{DSQRT}(0.999\text{D}0*\text{BIG})^{**2}$ in double precision) without overflowing. Similarly, `ETA` is the smallest floating-point number such that $\text{SQRT}(1.001*\text{ETA})^{**2}$ (or $\text{DSQRT}(1.001\text{D}0*\text{ETA})^{**2}$, respectively) does not underflow. `MACHEP` is the unit roundoff, that is, the smallest floating-point number such that $1 + \text{MACHEP}$ yields a stored floating-point number greater than 1. (Some computers feature registers that carry more bits than can be stored; `MACHEP` should only reflect the accuracy of numbers that can be stored.) `DATA` statements giving suitable values for `BIG`, `ETA`, and `MACHEP` for a variety of computers appear as comments in `RMDCON`.

Intrinsic functions are explicitly declared in the `NL2SOL` source code. On certain computers (e.g., Univac), it may be necessary to comment out these declarations. So that this may be done automatically by a simple program, such declarations are preceded by a comment having `C/+` in columns 1–3 and blanks in columns 4–72 and are followed by a comment having `C/` in columns 1 and 2 and blanks in columns 3–72.

3.13 Using Reverse Communication: `NL2ITR`

Instead of writing subroutines `CALCR` and `CALCJ` to compute the residual vector $R(X)$ and Jacobian matrix $J(X)$, one can call `NL2ITR` and provide `R` and `J` by reverse communication. The calling sequence is

```
CALL NL2ITR (D, IV, J, N, NN, P, R, V, X)
```

Parameters `IV`, `N`, `P`, `V`, and `X` are the same as the corresponding ones to `NL2SOL`, with the following exceptions: `V` need only contain $93 + 2N + P(3P + 31)/2$ elements, since the storage that `NL2SOL` and `NL2SNO` allocate for `D`, `J`, and `R` at the end of `V` is not needed; and components `IV(D)`, `IV(J)`, and `IV(R)` are not referenced. `D` is the scale vector (dimensioned `D(P)`). `NN` is the (integer) lead dimension for the `J` array, which is dimensioned `J(NN, P)`; `NN` must satisfy $NN \geq N$.

When `NL2ITR` is first called (with `IV(1) = 0` or `12`), `J` must have been set to $J(X)$, `R` to $R(X)$. When `NL2ITR` wants `R` to be evaluated at a new `X`, it returns with `IV(1) = 1`; the caller should then set `R` to $R(X)$ (unless `X` is out of range, in which case the caller should set `IV(TOOBIG)`, that is, `IV(2)`, to 1) and call `NL2ITR` again. Similarly, when `NL2ITR` wants `J` to be evaluated at `X`, it returns with `IV(1) = 2`, and the caller should then set `J` to $J(X)$ and call `NL2ITR` again. (If `J` cannot be evaluated at `X`, the caller may set `IV(NFGCAL)`, that is, `IV(7)`, to 0; this will cause `NL2ITR` to give the error return `IV(1) = 15`.)

3.14 `STOPX`

It is possible to arrange for `NL2SOL` (`NL2SNO` and `NL2ITR`) to be interrupted before each evaluation of $R(X)$ when used in an interactive environment. To do this, it is necessary to replace the logical function `STOPX` supplied with the `NL2SOL` package (which always returns `.FALSE.`) by a system-dependent `STOPX` that returns `.TRUE.` if and only if the “break” (i.e., “interrupt”) key has been pressed since the last call on `STOPX`. Once this is done, `NL2SOL` will return with `IV(1) = 11` when the “break” key is pressed before some other return has occurred. It is then possible to change some of the `IV` and `V` input components and restart; see Section 3.7.

3.15 Other V Input Values

V(COSMIN) . . . V(43) \in [MACHEP, 1] is the minimum absolute cosine allowed between the step just taken, Δx , and the corresponding change in gradients, Δg , for a full update of the S matrix to be made. If $|\Delta x^T \Delta g| / (\|\Delta x\|_2 \cdot \|\Delta g\|_2) < V(\text{COSMIN})$, then $\Delta x^T \Delta g$ is replaced in the update formula by $\text{sign}(\Delta x^T \Delta g) V(\text{COSMIN}) \|\Delta x\|_2 \|\Delta g\|_2$. Default = $\max\{10^{-6}, 100 \text{ MACHEP}\}$.

V(DECFACTOR) . . . V(22) \in [0.01, 0.8] is the factor by which the trust region radius is shrunk if CALCR sets NF to 0 (or NL2ITR is called with IV(1) = 1 and IV(TOOBIG) = 1). Default = 0.5.

V(EPSLON) . . . V(19) \in [0.001, 0.9] is the maximum relative difference allowed between $g^T \Delta x + 1/2 \Delta x^T H \Delta x$ and its optimal value subject to the constraint $\|\text{diag}(D) \Delta x\|_2 \leq V(\text{RADIUS})$, where Δx is the step being computed, g is the current gradient, and H is the current Hessian approximation. This is used in detecting and handling the special case discussed in [2]. Default = 0.1.

V(FUZZ) . . . V(45) \in [1.01, 100] is used in the test that decides whether to switch models. If q is the current model for F (near the point X) and \tilde{q} is the other model, and if

$$V(\text{FUZZ}) |\tilde{q}(X + \Delta x) - F(X + \Delta x)| < |q(X + \Delta x) - F(X + \Delta x)|,$$

then the models are switched. Default = 1.5.

V(INCFAC) . . . V(23) \in [1.2, 100] is the minimum factor by which the trust region radius is increased (when it is increased at all). Default = 2.

V(PHMNFC) . . . V(20) \in [-0.99, -0.001] is the minimum value allowed for $[\|\text{diag}(D) \Delta x\|_2 - V(\text{RADIUS})] / V(\text{RADIUS})$. Default = -0.1.

V(PHMXFC) . . . V(21) \in [1.2, 100] is the maximum value allowed for $[\|\text{diag}(D) \Delta x\|_2 - V(\text{RADIUS})] / V(\text{RADIUS})$. Default = 0.1.

V(RDFCMN) . . . V(24) \in [0.01, 0.8] is the minimum factor by which the trust region radius, V(RADIUS), may be shrunk. Default = 0.1.

V(RDFCMX) . . . V(25) \in [1.2, 100] is the maximum factor by which the trust region radius, V(RADIUS), may be increased at one time. Default = 4.0.

V(RLIMIT) . . . V(42) $\geq 10^{10}$ is the largest value allowed for $\|R(X)\|_2$ before F(X) is considered to overflow. Default = $(0.999 \cdot \text{BIG})^{1/2}$, where BIG is described in Section 3.12.

V(TUNER2) . . . V(27) \in [0, 0.5]. For a step to be accepted, the actual function reduction must be more than V(TUNER2) times its predicted value. Default = 0.0001.

V(TUNER3) . . . V(28) \in [0.001, 1]. If the actual function decrease is at least V(TUNER3) times the inner product of the step and the gradient (at the start of the step), then the trust region radius is increased. Default = 0.75.

V(TUNER4) . . . V(29) \in [-1, 1]. If the disposition of the new radius has not yet been decided and either

$$\|\text{diag}(D)^{-1} [H \Delta x - (g - g_0)]\| < V(\text{TUNER4}) \|\text{diag}(D)^{-1} g\|$$

or $g^T \Delta x < V(\text{TUNER5}) g_0^T \Delta x$, where Δx is the step just taken, H is the Hessian approximation used in computing Δx , g_0 is the old gradient, g is the new gradient, and D is the newly updated scale vector, then the radius is increased by a factor of V(INCFAC). Otherwise, it is left unchanged. Default = 0.5.

V(TUNER5) . . . V(30) $\geq \text{MACHEP}$ is described above with V(TUNER4). Default = 0.75.

3.16 Storage Requirements

NL2SOL, NL2SNO, and the subroutines from the NL2SOL package that they call amount to around 2360 FORTRAN statements (including nonexecutable statements, such as type statements, but excluding comments); the many comments bring the source code up to nearly 5200 lines. When compiled by the *H*-extended compiler on the IBM 370/168 at the Massachusetts Institute of Technology, this source code results in about 56,300 bytes of object code. The amount of variable storage used is listed above in Section 3.1.

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ALGORITHM

[The complete listing is available from the ACM Algorithms Distribution Service.]

ALGORITHM 574

Shape-Preserving Osculatory Quadratic Splines [E1, E2]

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and
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Key Words and Phrases: polynomial interpolation, osculation, shape preserving, convexity preserving, monotonicity preserving, Bernstein polynomial
CR Categories: 5.18, 8.2
Language: FORTRAN

DESCRIPTION

This algorithm is a FORTRAN implementation for the procedure developed in [3]. Let n data points $\{(x_i, y_i)\}_{i=1}^n$ and n first derivatives $\{m_i\}_{i=1}^n$ at these data points be given, with $x_i \leq x_{i+1}$, $1 \leq i \leq n-1$. The algorithm constructs a smooth osculatory quadratic spline S , which satisfies

- (1) $S(x_i) = y_i$, $1 \leq i \leq n$;
- (2) $S'(x_i) = m_i$, $1 \leq i \leq n$;
- (3) S preserves monotonicity and convexity in the case that the slopes m_i are consistent with the shape of the data;
- (4) the knots of the spline S include the data points and at most two additional knots between adjacent data points.

The spline S is a piecewise quadratic Bernstein polynomial with a continuous first derivative.

Subroutine SLOPES

This subroutine will calculate values for the slopes $\{m_i\}_{i=1}^n$, which will always be consistent with the shape of the data and which guarantee that at most one additional knot is required between adjacent data points. This subroutine is called before the main routine MEVAL if the user does not wish to provide these values.

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Subroutine MEVAL

The input to MEVAL is the set of data points $\{(x_i, y_i)\}_{i=1}^n$, the slopes $\{m_i\}_{i=1}^n$, and a vector of abscissas $u = (u_1, \dots, u_m)$ at which the spline S is to be evaluated. In addition, a nonnegative tolerance parameter EPS must be specified for use by the subroutine CHOOSE to distinguish various cases which the algorithm must consider. This parameter is described later. If an argument u_j is determined by the subroutine SEARCH to lie in the interval $[x_i, x_{i+1}]$, then MEVAL calls the subroutines CHOOSE and CASES, in that order, to compute the parameters of S for the interval. The subroutine then calls a function subprogram SPLINE which evaluates S at the argument u_j . Extrapolation is performed for those components of u that lie outside the interval $[x_1, x_n]$.

On exit from MEVAL the error parameter ERR is zero if processing proceeded normally and no extrapolation was required. ERR is 1 if at least one component of the vector u lies outside the interval $[x_1, x_n]$. A value of 2 indicates that the components of the vector u were not in ascending order, in which case no values of S are returned.

Subroutine SEARCH

This subroutine is a binary search used by MEVAL to locate the interval in which an argument of S lies.

Subroutine CHOOSE

The value of S on an interval $[x_i, x_{i+1}]$ depends only on the points (x_i, y_i) and (x_{i+1}, y_{i+1}) and the slopes m_i and m_{i+1} at these points. The subroutine determines which of Cases 1-4 applies [3]. The tolerance parameter EPS is used by CHOOSE to handle the pathologies that may occur. When the user provides values for the slopes other than those computed by the subroutine SLOPES, small perturbations may cause unwanted changes in the shape of the resulting spline. If the values of m_i or m_{i+1} are relatively close to either the slope S_i of the line joining the points (x_i, y_i) and (x_{i+1}, y_{i+1}) or to $2S_i$, roundoff may affect both the monotonicity and convexity of the spline. For example, if $|M - S_i| \leq \text{EPS} |S_i|$ for $M = m_i$ or m_{i+1} , then the algorithm assumes that $M = S_i$ and chooses either Case 2 or Case 3. We avoid Case 1 to preclude a nearly linear spline joining the two data points. Similarly, in order to avoid a spline with sharp bends at the endpoints, if $|m_i| > 2|S_i|$ and $|m_{i+1}| > (2 - \text{EPS})|S_i|$, then the algorithm assumes $|m_{i+1}| > 2|S_i|$ and selects Case 4 rather than Case 3. These choices are for aesthetics only and can easily be overridden by modifying the slopes m_i and m_{i+1} and/or setting EPS to a smaller value. If $\text{EPS} \neq 0$, then EPS should be greater than or equal to machine epsilon.

Subroutine CASES

After CHOOSE determines the correct case number, the subroutine CASES calculates the knots and other parameters which define S on the interval $[x_i, x_{i+1}]$. The spline S will have a single knot between x_i and x_{i+1} if NCASE is 1, 2, or 3. Case 4 requires two knots.

Function Subprogram SPLINE

Given the case number determined by CHOOSE and the parameters calculated by CASES, SPLINE determines the location of the argument relative to the knots in the interval and computes $S(u_j)$, which is the value of a quadratic Bernstein polynomial.

ACKNOWLEDGMENTS

The authors would like to thank Steven L. Dodd and Marie Roulier of North Carolina State University for their help in coding the algorithm.

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(Note. References [1, 2, 4, 5, 6, 7] are not cited in the text.)

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ALGORITHM

[The complete listing is available from the ACM Algorithms Distribution Service.]

ALGORITHM 575

Permutations for a Zero-Free Diagonal [F1]

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Key Words and Phrases: unsymmetric permutations, maximum transversal, maximum assignment, block triangular form, sparse matrices
CR Categories: 5.0, 5.1, 5.3, 5.4
Language: FORTRAN

DESCRIPTION

The subroutine is evoked by the FORTRAN statement
CALL MC21A(N, ICN, LICN, IP, LENR, IPERM, NUMNZ, IW)

where the parameters are described in the listing given here.

Given the pattern of nonzeros of a sparse matrix, this subroutine attempts to find a row permutation that makes the matrix have no zeros on its diagonal. It is possible that the user may input a matrix for which there is no permutation that makes the diagonal zero-free. An example of this is

$$\begin{pmatrix} X & 0 \\ X & 0 \end{pmatrix}$$

In such instances the algorithm will produce a permutation that will put as many nonzeros on the diagonal as possible (1 in the above example). This number will be output in NUMNZ. The array IPERM will still hold a permutation of the integers 1, 2, . . . , N, but in this case N-NUMNZ of the elements (IPERM(I), I) will be zero.

It is envisaged that a common use of the subroutine will be as the first part of a two-stage process for the determination of the block triangular form of a sparse matrix (see, for example, [1]). The second stage could be performed by Harwell subroutine MC13D [4], and the two stages are combined by using subroutine MC23A in the Harwell package MA28 [2].

The subroutine is based on a depth first search with look-ahead technique and is described in detail in [3]. There, numerical results from using this subroutine, which has been tested on a wide range of both structured and randomly generated matrices, are also discussed.

This subroutine has been written in American National Standard FORTRAN. Special comment cards have been included so that Harwell subroutine OE04A can be used to make an IBM FORTRAN version that uses half-length integers (INTEGER*2) for all the arrays except IP. This approximately halves the core requirements at the cost of restricting the order of the system to $2^{15} - 1$.

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ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```

C S IS THE STANDARD FORTRAN VERSION  SI/
C I IS THE IBM VERSION
C
      SUBROUTINE MC21A(N, ICN, LICN, IP, LENR, IPERM, NUMNZ, IW)
C
C DESCRIPTION OF PARAMETERS.
C INPUT VARIABLES  N, ICN, LICN, IP, LENR.
C OUTPUT VARIABLES IPERM, NUMNZ.
C
C N  ORDER OF MATRIX.
C ICN ARRAY CONTAINING THE COLUMN INDICES OF THE NON-ZEROS.  THOSE
C BELONGING TO A SINGLE ROW MUST BE CONTIGUOUS BUT THE ORDERING
C OF COLUMN INDICES WITHIN EACH ROW IS UNIMPORTANT AND WASTED
C SPACE BETWEEN ROWS IS PERMITTED.
C LICN LENGTH OF ARRAY ICN.
C IP  IP(I), I=1,2,...,N, IS THE POSITION IN ARRAY ICN OF THE FIRST
C COLUMN INDEX OF A NON-ZERO IN ROW I.
C LENR  LENR(I) IS THE NUMBER OF NON-ZEROS IN ROW I, I=1,2,...,N.
C IPERM CONTAINS PERMUTATION TO MAKE DIAGONAL HAVE THE SMALLEST
C NUMBER OF ZEROS ON IT.  ELEMENTS (IPERM(I),I) I=1, ... N ARE
C NON-ZERO AT THE END OF THE ALGORITHM UNLESS MATRIX
C IS STRUCTURALLY SINGULAR.  IN THIS CASE, (IPERM(I),I) WILL
C BE ZERO FOR N-NUMNZ ENTRIES.
C NUMNZ NUMBER OF NON-ZEROS ON DIAGONAL OF PERMUTED MATRIX.
C IW  WORK ARRAY ... SEE LATER COMMENTS.
C
      INTEGER IP(N)
C      INTEGER*2 ICN(LICN), LENR(N), IPERM(N), IW(N,4)  I/
      INTEGER ICN(LICN), LENR(N), IPERM(N), IW(N,4)
      CALL MC21B(N, ICN, LICN, IP, LENR, IPERM, NUMNZ, IW(1,1),
* IW(1,2), IW(1,3), IW(1,4))
      RETURN
      END
C/
      SUBROUTINE MC21B(N, ICN, LICN, IP, LENR, IPERM, NUMNZ, PR, ARP,
* CV, OUT)
      INTEGER IP(N)
C
C DIVISION OF WORK ARRAY IS NOW DESCRIBED.
C
C PR(I) IS THE PREVIOUS ROW TO I IN THE DEPTH FIRST SEARCH.
C ARP(I) IS ONE LESS THAN THE NUMBER OF NON-ZEROS IN ROW I
C WHICH HAVE NOT BEEN SCANNED WHEN LOOKING FOR A CHEAP ASSIGNMENT.
C CV(I) IS THE MOST RECENT ROW EXTENSION AT WHICH COLUMN I
C WAS VISITED.
C OUT(I) IS ONE LESS THAN THE NUMBER OF NON-ZEROS IN ROW I
C WHICH HAVE NOT BEEN SCANNED DURING ONE PASS THROUGH THE
C MAIN LOOP.
C
C      INTEGER*2 ICN(LICN), LENR(N), IPERM(N), PR(N), CV(N),  I/
C      1ARP(N), OUT(N)  I/
      INTEGER ICN(LICN), LENR(N), IPERM(N), PR(N), CV(N), ARP(N), OUT(N)
C
C      INITIALIZATION OF ARRAYS.
      DO 10 I=1,N
          ARP(I) = LENR(I) - 1
          CV(I) = 0
          IPERM(I) = 0
      10 CONTINUE

```



```

      NUMNZ = 0
C
C
C   MAIN LOOP.
C   EACH PASS ROUND THIS LOOP EITHER RESULTS IN A NEW ASSIGNMENT
C OR GIVES A ROW WITH NO ASSIGNMENT.
      DO 100 JORD=1,N
        J = JORD
        PR(J) = -1
        DO 70 K=1, JORD
C LOOK FOR A CHEAP ASSIGNMENT
          IN1 = ARP(J)
          IF (IN1.LT.0) GO TO 30
          IN2 = IP(J) + LENR(J) - 1
          IN1 = IN2 - IN1
          DO 20 II=IN1, IN2
            I = ICN(II)
            IF (IPERM(I).EQ.0) GO TO 80
20          CONTINUE
C   NO CHEAP ASSIGNMENT IN ROW.
          ARP(J) = -1
C   BEGIN LOOKING FOR ASSIGNMENT CHAIN STARTING WITH ROW J.
30          OUT(J) = LENR(J) - 1
C   INNER LOOP. EXTENDS CHAIN BY ONE OR BACKTRACKS.
          DO 60 KK=1, JORD
            IN1 = OUT(J)
            IF (IN1.LT.0) GO TO 50
            IN2 = IP(J) + LENR(J) - 1
            IN1 = IN2 - IN1
C   FORWARD SCAN.
            DO 40 II=IN1, IN2
              I = ICN(II)
              IF (CV(I).EQ.JORD) GO TO 40
C   COLUMN I HAS NOT YET BEEN ACCESSED DURING THIS PASS.
              J1 = J
              J = IPERM(I)
              CV(I) = JORD
              PR(J) = J1
              OUT(J1) = IN2 - II - 1
              GO TO 70
40          CONTINUE
C
C   BACKTRACKING STEP.
50          J = PR(J)
          IF (J.EQ.-1) GO TO 100
60          CONTINUE
C
70          CONTINUE
C
C   NEW ASSIGNMENT IS MADE.
80          IPERM(I) = J
          ARP(J) = IN2 - II - 1
          NUMNZ = NUMNZ + 1
          DO 90 K=1, JORD
            J = PR(J)
            IF (J.EQ.-1) GO TO 100
            II = IP(J) + LENR(J) - OUT(J) - 2
            I = ICN(II)
            IPERM(I) = J
90          CONTINUE
C
100         CONTINUE
C
C   IF MATRIX IS STRUCTURALLY SINGULAR, WE NOW COMPLETE THE
C PERMUTATION IPERM
          IF (NUMNZ.EQ.N) GO TO 150
          DO 110 I=1,N
            ARP(I) = 0
110         CONTINUE
          K = 0
          DO 130 I=1,N
            IF (IPERM(I).NE.0) GO TO 120
            K = K + 1
            OUT(K) = 1
            GO TO 130
120          J = IPERM(I)

```

```
      ARP(J) = 1
130 CONTINUE
      K = 0
      DO 140 I=1,N
        IF (ARP(I).NE.0) GO TO 140
        K = K + 1
        IOUK = OUT(K)
        IPERM(IOUK) = I
140 CONTINUE
150 RETURN
      END
```

ALGORITHM 576

A FORTRAN Program for Solving $\mathbf{Ax} = \mathbf{b}$

[F4]

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Key Words and Phrases: linear equations, Gaussian elimination, new pivoting strategy
CR Categories: 5.13, 5.14
Language: FORTRAN

DESCRIPTION

Given any $n \times n$ system of linear equations

$$\mathbf{Ax} = \mathbf{b} \quad (1)$$

and a number ϵ , the algorithm calculates the solution \mathbf{x} to (1), if $\epsilon \leq 0$, or an approximation solution $\mathbf{x}^{(k)}$ satisfying

$$\|\mathbf{b} - \mathbf{Ax}^{(k)}\|_{\infty} < \epsilon, \quad (2)$$

if $\epsilon > 0$. Furthermore, if \mathbf{A} appears to be singular, an approximate solution to (1) is still determined, although it does not satisfy (2).

The algorithm consists of Gaussian elimination combined with a new pivoting strategy based on the following fact. Suppose after the k th step ($k = 0, 1, \dots, n - 1$) of Gaussian elimination that $\mathbf{x}^{(k)} = [x_1^{(k)}, \dots, x_k^{(k)}, 0, \dots, 0]^T$ is the approximate solution obtained by back substitution in the first k equations. Then the right-hand sides of the last $n - k$ equations are the nonzero components of the residual vector $\mathbf{r}^{(k)} = \mathbf{b} - \mathbf{Ax}^{(k)}$; this is proved in [1]. Our pivoting strategy thus chooses the equation with the absolutely largest residual as the next equation to be eliminated, and the $(k + 1)$ th pivot is then selected as the absolutely largest of the $n - k$ coefficients on the left-hand side of the chosen equation. Finally, the pivot is positioned on the principal diagonal by, if necessary, a row interchange and a column interchange.

When $\epsilon > 0$, the algorithm terminates after k Gaussian elimination steps and yields $\mathbf{x}^{(k)}$ by back substitution in the first k equations. Normally, $\mathbf{x}^{(k)}$ has k nonzero components and it satisfies k of the eqs. (1) exactly (ignoring rounding errors). However, if ϵ is a *small* positive number, in order to satisfy (2) it may be necessary to determine the solution \mathbf{x} rather than $\mathbf{x}^{(k)}$: in such a case, k has the value n on output.

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The elimination is performed by subroutine MODGE, which terminates after stage k if there is not a suitably large residual among the remaining $(n - k)$ right-hand sides, or if the coefficient matrix of the remaining $(n - k)$ rows is considered to be singular. The first condition is detected by comparing all residuals with the real input parameter EPS. Hence EPS should be set to a small positive value representing an insignificant residual for cases in which premature termination is sought; otherwise EPS should be set less than or equal to zero. The second premature termination condition is controlled by the real input parameter TOLER. The equation with the absolutely largest residual is rejected as the pivot row if the absolute value of the pivot is not greater than TOLER. In this case complete pivoting is performed during the $(k + 1)$ th stage. However, if this produces no pivot greater in absolute value than TOLER, the elimination is terminated and an approximate solution is determined by back substitution in the first k equations.

The algorithm can thus be applied to any $n \times n$ system of linear equations, although it is particularly well suited to problems where the residuals (right-hand sides) can be made small by solving for fewer than n of the unknowns. Interpolation problems are often of this type, since some of the parameters x_j frequently turn out to be quite small in practice. Also, ill-conditioned (and even singular) systems of equations can often be handled satisfactorily by choosing appropriate values for EPS and TOLER (see Examples 3 and 4).

The accuracy of a computed solution to $\mathbf{Ax} = \mathbf{b}$ can usually be increased by iterative improvement, and subroutine REFINE is supplied for this purpose. A feature of this subroutine is that the $(t + 1)$ th residual vector \mathbf{r}_{t+1} is computed as

$$\mathbf{r}_{t+1} = \mathbf{r}_t - \mathbf{Ax}_t, \quad t = 0, 1, 2, \dots \quad (3)$$

where the residual \mathbf{r}_t and the computed solution \mathbf{x}_t are taken from the previous iteration, and $\mathbf{r}_0 = \mathbf{b}$. (It has been our experience that (3) permits better control of convergence than the more usual computation $\mathbf{r}_{t+1} = \mathbf{b} - \mathbf{A}(\mathbf{x}_0 + \mathbf{x}_1 + \dots + \mathbf{x}_t)$.) In this subroutine the output parameter DIGITS is used to estimate (usually fairly conservatively) the number of significant digits in the initial solution \mathbf{x}_0 . Basically, this is determined as $\log_{10} D$, where $D = \min_i |x_{0,i}/x_{1,i}|$, excluding values of i for which $x_{0,i} = 0$ or $x_{1,i} = 0$. However, if the quantity $\log_{10} D$ is greater than $\log_{10}(1/E)$, where $E = \max(\text{ERR}, 10^{-d})$ and d represents the number of decimal digits of accuracy available on the computer, then DIGITS is set equal to $\log_{10}(1/E)$ rather than $\log_{10} D$.

The following four examples were solved on an IBM 370/148 using a WATFIV compiler. All of the computations were performed in single-precision arithmetic (about 7 decimal digits) apart from the calculation of the residuals (3), which were formed in double-precision arithmetic and then assigned to a single-precision vector.

Example 1. Determine parameters x_1 , x_2 , and x_3 such that the quadratic polynomial $x_1 + x_2 t + x_3 t^2$ interpolates the data $\{(1, 89), (4, 209), (5, 201)\}$. Equivalently, solve the 3×3 system of linear equations

$$\begin{bmatrix} 1 & 1 & 1 \\ 1 & 4 & 16 \\ 1 & 5 & 25 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 89 \\ 209 \\ 201 \end{bmatrix}.$$

This example is used in [1] to illustrate our algorithm: the exact solution is $x_1 = 1$, $x_2 = 100$, and $x_3 = -12$.

Setting EPS = 0 and TOLER = 10^{-6} , subroutine MODGE gives $x_1 = 1.000010$, $x_2 = 99.999998$, and $x_3 = -11.999999$.

Setting EPS = 1 and TOLER = 10^{-6} , subroutine MODGE terminates after two Gaussian elimination steps and gives $x_1 = 0$, $x_2 = 100.4500$, and $x_3 = -12.04999$.

Example 2. Determine x_1, x_2, \dots, x_5 such that the function $x_1 t + x_2 t^2 + x_3 t^3 + x_4 t^4 + x_5 t^5$ interpolates te^{-t} for $t = 0.25(0.25)1.25$. Equivalently, solve $\mathbf{Ax} = \mathbf{b}$ where

$$\mathbf{A} = \begin{bmatrix} 0.25 & 0.25^2 & \dots & 0.25^5 \\ 0.50 & 0.50^2 & \dots & 0.50^5 \\ 0.75 & 0.75^2 & \dots & 0.75^5 \\ 1.00 & 1.00^2 & \dots & 1.00^5 \\ 1.25 & 1.25^2 & \dots & 1.25^5 \end{bmatrix} \quad \text{and} \quad \mathbf{b} = \begin{bmatrix} 0.25e^{-0.25} \\ 0.50e^{-0.50} \\ 0.75e^{-0.75} \\ 1.00e^{-1.00} \\ 1.25e^{-1.25} \end{bmatrix}.$$

Generating the elements of \mathbf{A} and \mathbf{b} within the computer (rather than reading them in as data), and setting $\text{EPS} = 0$ and $\text{TOLER} = 10^{-6}$, subroutine MODGE gives $x_1 = 0.9994718$, $x_2 = -0.9950786$, $x_3 = 0.4832436$, $x_4 = -0.1396533$, and $x_5 = 0.01989542$. (This solution contains almost one extra significant digit than the solution that can be obtained from Gaussian elimination with partial pivoting, using the modified version [3] of the subroutines in [2].) Iterative improvement is then applied to this initial solution.

Setting $\text{ERR} = 10^{-7}$ and $\text{ITER} = 14$, subroutine REFINE converges after two iterations to $x_1 = 0.9994708$, $x_2 = -0.9950703$, $x_3 = 0.4832247$, $x_4 = -0.1396357$, and $x_5 = 0.01988983$. The output parameter DIGITS indicates that the initial solution contains approximately 3.6 significant digits.

Example 3. Determine x_1, x_2, \dots, x_9 such that the function $x_1 + x_2t + x_3(t-1) + x_4t^2 + x_5(t^2-t) + x_6t^3 + x_7(t^3-t^2) + x_8t^4 + x_9(t^4-t^3)$ interpolates e^t for $t = 0(0.125)1$.

Setting $\text{EPS} = \text{TOLER} = 10^{-6}$, subroutine MODGE gives $x_1 = 0.9999993$, $x_2 = 0.9994056$, $x_3 = 0$, $x_4 = 0.5071687$, $x_5 = x_6 = x_7 = 0$, $x_8 = 0.2117074$, and $x_9 = -0.1439046$. The output parameters OCODE and PIVOT have the values 3 and 2, respectively, indicating that the computation was terminated prematurely because a pivot was chosen (by complete pivoting) whose magnitude is less than TOLER.

Setting $\text{ERR} = 10^{-7}$ and $\text{ITER} = 14$, subroutine REFINE converges after two iterations to $x_1 = 1.000000$, $x_2 = 0.9994109$, $x_4 = 0.5071452$, $x_8 = 0.2117255$, and $x_9 = -0.1439375$. DIGITS has the value 3.6 on output.

Example 4. Determine x_1, x_2, \dots, x_{32} such that the function

$$P(s, t) = x_1 + \sum_{j=1}^{15} (x_{2j} \text{Re}(z^j) + x_{2j+1} \text{Im}(z^j)) + x_{32} \text{Re}(z^{16}),$$

where $z = s + it$, interpolates $Q(s, t) = e^{s^2 - st + 2t^2}$ on the 32 points (s, t) for which $s = 0(0.125)1$, $t = 0(0.125)1$. This problem arises when attempting to solve approximately the Dirichlet problem for Laplace's equation on a unit square with boundary value Q . The coefficients x_j of the harmonic polynomial P are determined by collocation, that is, by requiring that $P = Q$ on the 32 specified boundary points. (We arranged the resulting 32 linear equations in order by proceeding in an anticlockwise direction around the boundary of the square, starting from the origin.)

Setting $\text{EPS} = 0$ and $\text{TOLER} = 10^{-6}$, subroutine MODGE gives an initial solution containing *no* significant digits: many of the 32 computed coefficients x_j have exponents of size 10^3 or 10^4 and only eight coefficients have magnitudes less than 20. (The initial solution that can be obtained from Gaussian elimination with partial pivoting also contains no significant digits.) Setting $\text{ERR} = 10^{-7}$ and $\text{ITER} = 14$, subroutine REFINE diverges, that is, all 14 iterations are performed without improvement in the accuracy of the coefficients x_j . The output parameter DIGITS has a *negative* value, clearly indicating that the initial solution contains no significant digits.

However, setting $\text{EPS} = 0.05$ and $\text{TOLER} = 10^{-6}$, subroutine MODGE terminates after 22 Gaussian elimination steps and gives an initial approximate solution in which 10 coefficients are set to zero and the remaining 22 coefficients vary from a minimum of $x_{27} = -9.698888$ to a maximum of $x_{29} = 12.78670$. Setting $\text{ERR} = 10^{-7}$ and $\text{ITER} = 14$, subroutine REFINE converges after three iterations: in this final solution $x_{27} = -9.699727$ and $x_{29} = 12.78701$. The output parameter

DIGITS indicates that the initial approximate solution contains approximately 2.9 significant digits.

This example illustrates how our algorithm can be used to produce a meaningful approximate solution to a system of linear equations whose numerical solution may be difficult to determine. The harmonic polynomial P^* , say, given by the above approximate solution satisfies $\|Q - P^*\|_\infty \leq 0.05$ on the 32 specified boundary bounds, and a fine grid search establishes that $\|Q - P^*\|_\infty \doteq 0.06$ on the complete boundary. By the maximum principle, it follows that the latter inequality also applies on the whole square.

ACKNOWLEDGMENTS

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(Note. References [4, 5] are not cited in the text.)

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ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```

      SUBROUTINE MODGE(N, NDIM, A, B, EPS, TOLER, P, Q, Z,
* X, K, OCODE, OPIVOT)
C THIS SUBROUTINE SOLVES ANY N-BY-N SYSTEM OF LINEAR
C EQUATIONS AX=B, USING A MODIFICATION OF GAUSSIAN
C ELIMINATION. THE MODIFICATION OCCURS IN THE CHOICE OF
C PIVOTS: AT EACH STAGE THE PIVOT ROW IS DETERMINED BY THE
C ABSOLUTELY LARGEST RESIDUAL (RIGHT HAND SIDE), AND THEN
C THE PIVOT IS CHOSEN AS THE ABSOLUTELY LARGEST COEFFICIENT
C IN THE PIVOT ROW. THE PIVOT IS POSITIONED ON THE PRINCIPAL
C DIAGONAL BY (IF NECESSARY) A ROW INTERCHANGE AND A COLUMN
C INTERCHANGE.
C
C IF AFTER THE K-TH STAGE THE LARGEST RESIDUAL IS TOO SMALL
C (.LT.EPS), THE ELIMINATION IS TERMINATED. BACK
C SUBSTITUTION IS THEN USED TO OBTAIN AN APPROXIMATE
C SOLUTION WHICH SATISFIES THE FIRST K EQUATIONS.
C
C IF THE (K+1)-TH PIVOT IS TOO SMALL (.LE.TOLER), THE
C ELIMINATION IS TERMINATED. BACK SUBSTITUTION IS THEN USED
C TO OBTAIN AN APPROXIMATE SOLUTION WHICH SATISFIES THE
C FIRST K EQUATIONS.
C
C DESCRIPTION OF PARAMETERS:
C N      AN INTEGER DENOTING THE NUMBER OF EQUATIONS
C        (N. GE. 1).
C NDIM   AN INTEGER PARAMETER FOR ADJUSTABLE DIMENSIONS
C        (NDIM. GE. N).
C A      A REAL ARRAY OF NDIM ROWS AND AT LEAST N
C        COLUMNS WHICH ON ENTRY MUST CONTAIN THE COEFFICIENT
C        MATRIX A IN ITS FIRST N ROWS AND N COLUMNS.
C        ON EXIT THE FIRST K ROWS AND K COLUMNS (K. LE. N) OF
C        A CONTAIN THE LU DECOMPOSITION OF THE K EQUATIONS
C        SATISFIED.
C B      A REAL ARRAY OF DIMENSION AT LEAST N WHICH

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C      ON ENTRY MUST CONTAIN IN POSITIONS 1 TO N THE RIGHT
C      HAND SIDE VECTOR B.
C      ON EXIT THE POSITIONS K+1 TO N CONTAIN THE RESIDUALS
C      OF THE UNSATISFIED EQUATIONS.
C EPS   A REAL INPUT TOLERANCE. IF THE (ABSOLUTELY) LARGEST
C      RESIDUAL AT ANY STAGE OF THE ELIMINATION IS LESS
C      THAN EPS THEN THE ELIMINATION IS TERMINATED AND
C      BACK SUBSTITUTION COMMENCES. PREMATURE TERMINATION
C      OF THIS TYPE CAN BE AVOIDED BY SETTING EPS TO ZERO.
C TOLER A REAL POSITIVE INPUT TOLERANCE. IF THE ABSOLUTE
C      VALUE OF THE PIVOT AT ANY STAGE OF THE ELIMINATION
C      IS LESS THAN OR EQUAL TO TOLER, THEN COMPLETE
C      PIVOTING IS USED FOR THIS STAGE. IF THE ABSOLUTE
C      VALUE OF THE NEW PIVOT IS LESS THAN OR EQUAL TO
C      TOLER, THE REMAINING EQUATIONS ARE CONSIDERED
C      TO BE LINEARLY DEPENDENT, AND SO THE ELIMINATION
C      IS TERMINATED AND BACK SUBSTITUTION COMMENCES.
C      TOLER SHOULD NORMALLY BE SET TO APPROXIMATELY
C      10**(-D+1), WHERE D REPRESENTS THE NUMBER OF
C      DECIMAL DIGITS OF ACCURACY AVAILABLE. HOWEVER,
C      PROBLEMS WHICH ARE NOT 'REASONABLY SCALED'
C      BEFOREHAND COULD REQUIRE A MUCH LARGER OR
C      SMALLER VALUE FOR TOLER THAN IS RECOMMENDED
C      HERE.
C P     AN INTEGER ARRAY OF DIMENSION AT LEAST N USED TO
C      RECORD COLUMN INTERCHANGES OF THE ARRAY A.
C      ON EXIT P(I) IS THE INDEX OF THE ORIGINAL POSITION
C      OF COLUMN I.
C Q     AN INTEGER ARRAY OF DIMENSION AT LEAST N USED TO
C      RECORD ROW INTERCHANGES OF THE ARRAYS A AND B.
C      ON EXIT Q(I) IS THE INDEX OF THE ORIGINAL POSITION
C      OF ROW I.
C Z     A REAL WORK SPACE ARRAY OF DIMENSION AT LEAST N.
C X     A REAL OUTPUT ARRAY OF DIMENSION AT LEAST N
C      WHICH CONTAINS THE SOLUTION (IN CORRECT ORDER) OF
C      THE N LINEAR EQUATIONS IN POSITIONS 1 TO N.
C      IF THE SOLUTION IS OBTAINED USING LESS THAN N
C      EQUATIONS, THEN THE REMAINING COMPONENTS OF X ARE
C      SET TO ZERO.
C K     AN INTEGER WHICH ON OUTPUT DENOTES THE NUMBER OF
C      EQUATIONS SATISFIED (K.LE.N).
C OCODE AN INTEGER EXIT CODE WITH VALUES:
C      0- EITHER N.LE.0 OR (N.EQ.1 AND A(1,1).EQ.0).
C      1- THE SYSTEM AX=B IS SOLVED.
C      2- THE ELIMINATION IS TERMINATED BY USE OF
C      EPS (SMALL RESIDUAL).
C      3- THE ELIMINATION IS TERMINATED BY USE OF
C      TOLER (SMALL PIVOT).
C OPIVOT AN INTEGER EXIT CODE WITH VALUES:
C      1- COMPLETE PIVOTING IS NOT EMPLOYED.
C      2- COMPLETE PIVOTING IS EMPLOYED, HENCE SMALL
C      RESIDUALS (.LT.EPS) MAY HAVE BEEN ALLOWED.
C
C AN IMPLEMENTATION NOTE:
C IF YOUR FORTRAN COMPILER PERMITS A SINGLE COLUMN OF A TWO
C DIMENSIONAL ARRAY TO BE PASSED TO A ONE DIMENSIONAL ARRAY
C THROUGH A SUBROUTINE CALL, CONSIDERABLE SAVINGS IN
C EXECUTION TIME MAY BE ACHIEVED BY USING SUBROUTINE
C SAXPY (SEE A.C.M. TRANSACTIONS ON MATHEMATICAL
C SOFTWARE VOL.5 NUM.3 1979 PP.308-323)
C      SUBROUTINE REFINE(N, NDIM, A, B, ALU, P, Q, KMAX,
C      * ERR, R, S, XNEW, DX, X, ITER, DIGITS)
C THIS SUBROUTINE PERFORMS ITERATIVE IMPROVEMENT ON THE
C COMPUTED SOLUTION X OBTAINED FROM SUBROUTINE MODGE.
C THE ITERATIONS CONTINUE UNTIL EITHER THE CORRECTIONS
C TO X ARE NEGLIGIBLE, I.E. CONVERGENCE OCCURS (SEE
C DESCRIPTION OF ERR), OR UNTIL COMPLETION OF THE
C MAXIMUM NUMBER OF ITERATIONS ALLOWED (ITER).
C
C THE FOLLOWING TABLE GIVES THE CORRESPONDENCE BETWEEN THE
C OUTPUT PARAMETERS OF MODGE AND THE INPUT PARAMETERS OF

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C REFINE:
C NAME ON OUTPUT FROM MODGE: NAME ON INPUT TO REFINE:
C N N
C NDIM NDIM
C A ALU
C P P
C Q Q
C K KMAX
C X X
C
C DESCRIPTION OF PARAMETERS:
C N AN INTEGER DENOTING THE NUMBER OF EQUATIONS
C (N. GE. 1).
C NDIM AN INTEGER PARAMETER FOR ADJUSTABLE DIMENSIONS
C (NDIM. GE. N).
C A A REAL ARRAY OF NDIM ROWS AND AT LEAST N COLUMNS
C WHICH ON ENTRY MUST CONTAIN THE COEFFICIENT
C MATRIX A IN ITS FIRST N ROWS AND N COLUMNS.
C A IS NOT ALTERED BY REFINE.
C B A REAL ARRAY OF DIMENSION AT LEAST N WHICH ON
C ENTRY MUST CONTAIN IN POSITIONS 1 TO N THE RIGHT
C HAND SIDE VECTOR B. B IS NOT ALTERED BY REFINE.
C ALU THE REAL OUTPUT ARRAY A OF SUBROUTINE MODGE. ALU IS
C NOT ALTERED BY REFINE.
C P AN INTEGER ARRAY OF DIMENSION AT LEAST N, USED
C TO RECORD COLUMN INTERCHANGES OF THE ARRAY A.
C P IS THE OUTPUT ARRAY P OF SUBROUTINE MODGE.
C Q AN INTEGER ARRAY OF DIMENSION AT LEAST N, USED
C TO RECORD ROW INTERCHANGES OF THE ARRAYS A AND B.
C Q IS THE OUTPUT ARRAY Q OF SUBROUTINE MODGE.
C KMAX THE INTEGER OUTPUT PARAMETER K OF MODGE.
C ERR A REAL INPUT PARAMETER USED TO TEST FOR
C CONVERGENCE. ERR SHOULD NOT BE LESS THAN 10**(-D),
C WHERE D REPRESENTS THE NUMBER OF DECIMAL DIGITS
C OF ACCURACY AVAILABLE.
C THE CONVERGENCE CRITERION IS:
C EXIT IF THE FOLLOWING RELATIONSHIP IS TRUE FOR
C ALL I=1 TO KMAX:
C ABS(DX(I)) .LE. ABS(ERR*X(I)). OR. X(I) .EQ. 0.0 .
C R A REAL WORK SPACE ARRAY OF DIMENSION AT LEAST N.
C (USED BY REFINE TO CALCULATE THE RESIDUALS AT EACH
C ITERATION.)
C S A REAL WORK SPACE ARRAY OF DIMENSION AT LEAST N.
C (USED BY REFINE TO CALCULATE THE CORRECTIONS, DX,
C AT EACH ITERATION.)
C XNEW A REAL WORK SPACE ARRAY OF THE SAME DIMENSION AS X.
C DX A REAL WORK SPACE ARRAY OF DIMENSION AT LEAST N.
C (USED BY REFINE TO CALCULATE A CORRECTION TO X
C AT EACH ITERATION.)
C X THE REAL OUTPUT ARRAY X OF SUBROUTINE MODGE.
C ON EXIT X CONTAINS (USUALLY) A BETTER SOLUTION
C TO THE LINEAR SYSTEM AX=B.
C ITER AN INTEGER DENOTING, ON INPUT, THE MAXIMUM NUMBER
C OF ITERATIONS OF ITERATIVE IMPROVEMENT ALLOWED.
C ITER SHOULD NOT EXCEED 2*D, WHERE D REPRESENTS
C THE NUMBER OF DECIMAL DIGITS OF ACCURACY AVAILABLE.
C ON OUTPUT ITER IS THE ACTUAL NUMBER OF ITERATIONS
C COMPLETED.
C DIGITS A REAL PARAMETER WHICH ON OUTPUT INDICATES THE
C APPROXIMATE NUMBER OF SIGNIFICANT DIGITS IN THE
C INITIAL SOLUTION (I.E. THE SOLUTION FROM
C MODGE).

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ALGORITHM 577

Algorithms for Incomplete Elliptic Integrals [S21]

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Key Words and Phrases: elliptic integrals, logarithms, inverse circular functions, inverse hyperbolic functions, R -functions
CR Categories: 5.12
Language: FORTRAN

DESCRIPTION

Robust and highly portable FORTRAN programs, implementing algorithms proved in [1], are presented for the symmetric elliptic integral of the first kind,

$$R_F(x, y, z) = \frac{1}{2} \int_0^{\infty} [(t+x)(t+y)(t+z)]^{-1/2} dt, \quad (1)$$

and a symmetric elliptic integral of the third kind,

$$R_J(x, y, z, \rho) = \frac{3}{2} \int_0^{\infty} [(t+x)(t+y)(t+z)]^{-1/2} (t+\rho)^{-1} dt. \quad (2)$$

All arguments are nonnegative, and if $x = 0$, the integrals are complete. The accessible range of argument values can be extended beyond the range admissible in the programs by using homogeneity:

$$R_F(kx, ky, kz) = k^{-1/2} R_F(x, y, z), \quad R_J(kx, ky, kz, k\rho) = k^{-3/2} R_J(x, y, z, \rho).$$

The range of function values is indicated by

$$\begin{aligned} M^{-1/2} \leq R_F(x, y, z) < 2m^{-1/2}, & \quad M_1^{-3/2} \leq R_J(x, y, z, \rho) < 5m_1^{-3/2}, \\ M = \max\{x, y, z\}, & \quad M_1 = \max\{M, \rho\}, \\ m = \min\{x+y, x+z, y+z\}, & \quad m_1 = \min\{m, \rho\}. \end{aligned}$$

Enough simplifications occur in degenerate cases to justify separate programs for

$$R_C(x, y) = R_F(x, y, y) \quad \text{and} \quad R_D(x, y, z) = R_J(x, y, z, z). \quad (3)$$

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The latter is an integral of the second kind, while R_C is an elementary function from which inverse circular functions can be computed if $x < y$ and logarithms or inverse hyperbolic functions if $x > y$ (see [1] for explicit formulas). The subroutine RJ calls RC in each cycle of iteration. All four subroutines proceed by successive applications of the duplication theorem, followed by expansion in Taylor series to fifth order. The programs require little space, convergence is linear but quite fast, serious cancellations do not occur, and accuracy is adjustable.

Legendre's elliptic integrals can be computed from these functions, but enumeration of cases is avoided (because of symmetry) by expressing a given elliptic integral directly in terms of R_F , R_D , and R_J . Compact tables of elliptic integrals in terms of these functions are being prepared (e.g., see [2]).

The codes have been verified by PFORT, but fall short of complete portability because they involve two machine-dependent constants, LOLIM and UPLIM, which define the range of admissible arguments. No argument may be greater than UPLIM. The role of LOLIM is less simple, but in most cases it is a lower bound for the sum of two arguments. For R_C and R_F , which are homogeneous of degree $-\frac{1}{2}$, LOLIM and UPLIM differ from the machine minimum and maximum, respectively, by a factor of only 5. For R_D and R_J , which are homogeneous of degree $-\frac{3}{2}$, the limits are inevitably much more restrictive. Values of LOLIM and UPLIM are suggested in the program comments for several types of computers, and formulas are given for use with other machines. A larger (smaller) value of LOLIM (UPLIM) may be used with no ill effects, except that the range of valid arguments will be narrowed. Since the subroutine RJ calls RC, the values of LOLIM and UPLIM for the two programs must be coordinated as specified in the comments.

Our codes have been modified by Dr. J.L. Schonfelder for inclusion in a future edition of the NAG library, and in the case of RC and RF his versions admit all machine-representable arguments. He uses LOLIM and UPLIM to define various regions in which different scaling procedures are used. This entails a significant elaboration of the codes, but will be valuable for users who need to work near the fringes of machine representability.

The accuracy of each program is controlled by a variable called ERRTOL, which determines the maximum error committed in truncating a Taylor series at terms of fifth order. If lower precision is sufficient, a larger value of ERRTOL will speed up the program by requiring fewer iterations of the duplication theorem before computing the Taylor series. Decreasing ERRTOL by a factor of 10 yields six more decimal digits of accuracy at the expense of one or two more iterations. More precisely, decreasing ERRTOL by a factor of 4 requires one more iteration and reduces the error ultimately by a factor of $4^6 = 4096$. Since required accuracy may vary from one run to another, ERRTOL is treated as an input variable. (Its value is set automatically when RC is called by RJ.) On the other hand, LOLIM and UPLIM are set within each subroutine because they will usually be fixed once and for all at a given installation. An exception to the last statement occurs because UPLIM in the case of RD depends on ERRTOL (essentially because the number of iterations depends on ERRTOL, and underflow in computing SIGMA becomes more likely in successive iterations as POWER4 decreases). For nearly all purposes ERRTOL may be replaced by 10^{-5} in the formula for UPLIM, but users who need to admit very large arguments or who want 30 decimal digits of accuracy without losing robustness should recompute UPLIM from the value actually used for ERRTOL.

The programs have been tested for robustness on an ITEL AS/6 computer by running them at extreme points of the region of valid arguments (such as $X = 0$ and $Y = \text{UPLIM}$ in the case of RC). In order to achieve robustness the Taylor series given in [1] had to be rearranged in terms of elementary symmetric functions because underflows occurred in computing the power sums $s_n^{(m)}$. By means of [1, e.g., (5.10)] the series in [1, eq. (2.5)] becomes

$$1 - \frac{1}{10}E_2 + \frac{1}{14}E_3 + \frac{1}{24}E_2^2 - \frac{3}{44}E_2E_3 + r_n, \quad (4)$$

where E_r is the elementary symmetric function (see [1, Appendix]) of degree r in X_n, Y_n, Z_n . Similarly, the series in [1, (2.19) and (2.27)] become

$$1 - \frac{3}{14}E_2 + \frac{1}{6}E_3 - \frac{3}{22}E_4 + \frac{9}{88}E_2^2 + \frac{3}{26}E_5 - \frac{9}{52}E_2E_3 + r_n, \quad (5)$$

where E_r is an elementary symmetric function of X_n, Y_n, Z_n, P_n, P_n for the case of [1, (2.19)] and X_n, Y_n, Z_n, Z_n, Z_n for [1, (2.27)]. In all cases, $E_1 = 0$.

Some check values are

$$\begin{aligned} R_C(0, \frac{1}{4}) &= R_C(\frac{1}{16}, \frac{1}{8}) = \pi, \\ R_C(\frac{9}{4}, 2) &= \log_e 2, \\ R_F(0, 1, 2) &= A = 1.31102\ 87771\ 46059\ 90523\ 24198, \\ R_D(0, 2, 1) &= 3B = 3\pi/4A = 1.79721\ 03521\ 03388\ 31115\ 98837, \\ R_J(2, 3, 4, 5) &= 0.14297\ 57966\ 71567\ 53833\ 23308. \end{aligned} \quad (6)$$

The lemniscate constants A and B are given to 50D in [3]. Consistency checks are provided in the program comments. Values of $(\log x)/(x - 1)$ and $\arctan x$ computed from library routines for the AS/6 computer agree well with those computed from RC, usually to within four units in the last (sixteenth) decimal place. There is similar agreement between RF and the FUNPACK routine for the complete elliptic integral $K(k)$. Because of homogeneity a substantially wider range of values of k^2 can be used with RF than with FUNPACK.

ACKNOWLEDGMENTS

We are grateful for helpful comments and suggestions from J.L. Schonfelder and Edward W. Ng. An important misprint was pointed out by Wayne Fullerton. We owe an exceptional debt to Robert L. Pexton for his interest, encouragement, and unselfish help. He ran many tests on a CDC 7600, and his excellent suggestions and advice have made the programs far better than they would otherwise have been.

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2. CARLSON, B.C. Elliptic integrals of the first kind. *SIAM J. Math. Anal.* 8 (1977), 231-242.
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ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```

C          *****
C
C          DOUBLE PRECISION FUNCTION RC(X, Y, ERRTOL, IERR)
C
C          THIS FUNCTION SUBROUTINE COMPUTES THE ELEMENTARY INTEGRAL
C          RC(X, Y) = INTEGRAL FROM ZERO TO INFINITY OF
C
C
C
C
C          (1/2)(T+X)  -1/2  (T+Y)  -1  DT,
C
C          WHERE X IS NONNEGATIVE AND Y IS POSITIVE.  THE DUPLICATION
C          THEOREM IS ITERATED UNTIL THE VARIABLES ARE NEARLY EQUAL,
C          AND THE FUNCTION IS THEN EXPANDED IN TAYLOR SERIES TO FIFTH
C          ORDER.  LOGARITHMIC, INVERSE CIRCULAR, AND INVERSE HYPER-
C          BOLIC FUNCTIONS CAN BE EXPRESSED IN TERMS OF RC.  REFERENCE:
C          B. C. CARLSON, COMPUTING ELLIPTIC INTEGRALS BY DUPLICATION,
C          NUMER. MATH. 33 (1979), 1-16.  CODED BY B. C. CARLSON AND
C          ELAINE M. NOTIS, AMES LABORATORY-DOE, IOWA STATE UNIVERSITY,
C          AMES, IOWA 50011.  MARCH 1, 1980.
C
C          CHECK BY ADDITION THEOREM: RC(X, X+Z) + RC(Y, Y+Z) = RC(0, Z),
C          WHERE X, Y, AND Z ARE POSITIVE AND X * Y = Z * Z.

```

```

C
  INTEGER IERR, PRINTR
  DOUBLE PRECISION C1, C2, ERRTOL, LAMDA, LOLIM
  DOUBLE PRECISION MU, S, SN, UPLIM, X, XN, Y, YN
C
  INTRINSIC FUNCTIONS USED: DABS, DMAX1, DSQRT
C
  PRINTR IS THE UNIT NUMBER OF THE PRINTER.
  DATA PRINTR /6/
C
  LOLIM AND UPLIM DETERMINE THE RANGE OF VALID ARGUMENTS.
  LOLIM IS NOT LESS THAN THE MACHINE MINIMUM MULTIPLIED BY 5.
  UPLIM IS NOT GREATER THAN THE MACHINE MAXIMUM DIVIDED BY 5.
C
  ACCEPTABLE VALUES FOR:      LOLIM      UPLIM
C
  IBM 360/370 SERIES      :   3.D-78      1.D+75
C
  CDC 6000/7000 SERIES    :   1.D-292     1.D+321
C
  UNIVAC 1100 SERIES      :   1.D-307     1.D+307
C
  WARNING: IF THIS PROGRAM IS CONVERTED TO SINGLE PRECISION,
  THE VALUES FOR THE UNIVAC 1100 SERIES SHOULD BE CHANGED TO
  LOLIM = 1.E-37 AND UPLIM = 1.E+37 BECAUSE THE MACHINE
  EXTREMA CHANGE WITH THE PRECISION.
C
  DATA LOLIM /3.D-78/, UPLIM /1.D+75/
C
  ON INPUT:
C
  X AND Y ARE THE VARIABLES IN THE INTEGRAL RC(X,Y).
C
  ERRTOL IS SET TO THE DESIRED ERROR TOLERANCE.
  RELATIVE ERROR DUE TO TRUNCATION IS LESS THAN
  16 * ERRTOL ** 6 / (1 - 2 * ERRTOL).
C
  SAMPLE CHOICES:  ERRTOL      RELATIVE TRUNCATION
C
C                   1.D-3      ERROR LESS THAN
C                   3.D-3      2.D-17
C                   1.D-2      2.D-14
C                   3.D-2      2.D-11
C                   1.D-1      2.D-8
C                   1.D-1      2.D-5
C
  ON OUTPUT:
C
  X, Y, AND ERRTOL ARE UNALTERED.
C
  IERR IS THE RETURN ERROR CODE:
  IERR = 0 FOR NORMAL COMPLETION OF THE SUBROUTINE.
  IERR = 1 FOR ABNORMAL TERMINATION.
C
  *****
  WARNING: CHANGES IN THE PROGRAM MAY IMPROVE SPEED AT THE
  EXPENSE OF ROBUSTNESS.
C
  IF (X.LT.0.DO.OR.Y.LE.0.DO) GO TO 10
  IF ((X+Y).LT.LOLIM) GO TO 10
  IF (DMAX1(X,Y).LE.UPLIM) GO TO 20
10 WRITE (PRINTR,99999)
  WRITE (PRINTR,99998) X, Y
  IERR = 1
  GO TO 50
C
20 IERR = 0
  XN = X
  YN = Y
C
30 MU = (XN+YN+YN)/3.DO
  SN = (YN+MU)/MU - 2.DO
  IF (DABS(SN).LT.ERRTOL) GO TO 40
  LAMDA = 2.DO*DSQRT(XN)*DSQRT(YN) + YN
  XN = (XN+LAMDA)*0.25DO
  YN = (YN+LAMDA)*0.25DO
  GO TO 30
C
40 C1 = 1.DO/7.DO
  C2 = 9.DO/22.DO
  S = SN*SN*(0.3DO+SN*(C1+SN*(0.375DO+SN*C2)))
  RC = (1.DO+S)/DSQRT(MU)

```


ALGORITHM 578

Solution of Real Linear Equations in a Paged Virtual Store [F4]

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Key Words and Phrases: Gaussian elimination, paged virtual store

CR Categories: 5.14

Language: FORTRAN

DESCRIPTION

BLCFAC is a routine for performing Gaussian elimination with partial pivoting on a real square matrix A , with the operations on blocks of consecutive columns grouped together to minimize the number of page swaps on a machine with a paged virtual store, as described in [1]. Given that A has been factorized by BLCFAC, the routine BLCSOL will solve either of the systems

$$AX = B$$

or

$$A^T X = B$$

with multiple right-hand sides. The operations on blocks of consecutive right-hand sides are grouped together for the same reason. BLCFAC and BLCSOL are essentially the same as the routines F01BTF and F04AYF in the NAG FORTRAN Library.

The routines need two environmental parameters, which are supplied in the functions SRELPR and NSACT. SRELPR is the relative precision, as defined in [2]; NSACT is the target active set size in real storage units. The choice of a suitable value, which depends on the particular machine, is discussed in Section 4 of [1]. The value zero is suitable on nonpaged machines.

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DOCUMENTATION

The routines are fully documented by embedded comments.

EXAMPLE

To solve the set of linear equations $Ax = b$, where

$$A = \begin{bmatrix} 3 & -6 & 12 & 3 \\ -1 & 4 & -10 & -1 \\ 4 & -13 & 30 & 2 \\ 2 & -11 & 30 & 7 \end{bmatrix} \quad \text{and} \quad b = \begin{bmatrix} -3 \\ 1 \\ -9 \\ 9 \end{bmatrix}$$

```

INTEGER I, IA, IB, IFAIL, IR, J, N, NIN, NOUT
LOGICAL TRANS
REAL A(5, 5), B(5), P(5)
DATA NIN /5/, NOUT /6/
N = 4
READ (NIN, 99999) ((A(I, J), J = 1, N), I = 1, N)
IA = 5
IFAIL = 1
C   GAUSSIAN ELIMINATION
CALL BLCFAC(N, A, IA, P, IFAIL)
IF (IFAIL.EQ.0) GO TO 20
WRITE (NOUT, 99998) IFAIL
STOP
20  IR = 1
    TRANS = .FALSE.
    READ (NIN, 99999) (B(I), I = 1, N)
    IB = 5
C   APPROXIMATE SOLUTION OF LINEAR EQUATIONS
CALL BLCSOL(N, IR, TRANS, A, IA, P, B, IB, IFAIL)
WRITE (NOUT, 99997) (B(I), I = 1, N)
STOP
99999 FORMAT (4F5.0)
99998 FORMAT (25H0ERROR IN BLCFAC IFAIL = , I2)
99997 FORMAT (10H0SOLUTIONS/(1H , F4.1))
END

```

The following are suitable data:

```

      3  -6  12  3
     -1   4 -10 -1
      4 -13  30  2
      2 -11  30  7
     -3   1  -9  9

```

and the following results should be obtained:

SOLUTIONS

```

     -1.0
      3.0
      1.0
      2.0

```

REFERENCES

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2. FORD, B. Parameterization of the environment for transportable numerical software. *ACM Trans. Math. Softw.* 4, 2 (June 1978), 100-103.

ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```

SUBROUTINE BLCFAC(N, A, IA, P, IERR)                                BLF 10
C                                                                    BLF 20
C BLCFAC FACTORIZES A REAL MATRIX BY GAUSSIAN ELIMINATION WITH    BLF 30
C PARTIAL PIVOTING . IT IS ORGANIZED SO THAT BLOCKS OF           BLF 40
C CONSECUTIVE COLUMNS ARE TREATED TOGETHER FOR EFFICIENCY ON A   BLF 50
C PAGED MACHINE.                                                BLF 60
C                                                                    BLF 70
C THE ROUTINE FACTORIZES THE INPUT MATRIX A AS P * L * U,        BLF 80
C WHERE P IS A PERMUTATION MATRIX, L IS A LOWER TRIANGULAR       BLF 90
C MATRIX , AND U IS AN UPPER TRIANGULAR MATRIX WITH UNIT         BLF 100
C DIAGONAL ELEMENTS. WHEN CHOOSING A PIVOT, THE ROUTINE          BLF 110
C IMPLICITLY SCALES THE ROWS TO HAVE LARGEST ELEMENT 1.0. THE    BLF 120
C ROUTINE TESTS THE SIZE OF THE IMPLICITLY SCALED PIVOT ELEMENT  BLF 130
C TO CHECK FOR APPROXIMATE SINGULARITY.                          BLF 140
C                                                                    BLF 150
C PARAMETERS-                                                    BLF 160
C                                                                    BLF 170
C N      INTEGER                                                BLF 180
C ON ENTRY, N SPECIFIES THE ORDER OF THE MATRIX A                BLF 190
C (N.GT.0). UNCHANGED ON EXIT.                                  BLF 200
C                                                                    BLF 210
C A      REAL ARRAY OF DIMENSION (IA,N)                          BLF 220
C BEFORE ENTRY, A(I,J) MUST BE SET TO THE (I,J)TH ELEMENT       BLF 230
C OF THE MATRIX A, FOR I = 1,...,N AND J = 1,...,N.              BLF 240
C ON SUCCESSFUL EXIT, THE STRICT UPPER TRIANGLE OF A IS         BLF 250
C OVERWRITTEN BY THE CORRESPONDING ELEMENTS OF U (THE           BLF 260
C UNIT DIAGONAL ELEMENTS OF U ARE NOT STORED) AND THE           BLF 270
C LOWER TRIANGLE OF A IS OVERWRITTEN BY THE MULTIPLIERS         BLF 280
C USED IN THE GAUSSIAN ELIMINATION (IT DOES NOT CONTAIN         BLF 290
C THE MATRIX L AS SUCH - THE SUBDIAGONAL ELEMENTS ARE          BLF 300
C PERMUTED).                                                    BLF 310
C                                                                    BLF 320
C IA     INTEGER                                                BLF 330
C ON ENTRY, IA SPECIFIES THE FIRST DIMENSION OF A AS            BLF 340
C DECLARED IN THE CALLING (SUB)PROGRAM (IA.GE.N).               BLF 350
C UNCHANGED ON EXIT.                                           BLF 360
C                                                                    BLF 370
C P      REAL ARRAY OF DIMENSION (N)                             BLF 380
C ON SUCCESSFUL EXIT, P(I) CONTAINS THE ROW INDEX OF            BLF 390
C THE I-TH PIVOT, FOR I = 1,...,N.                              BLF 400
C                                                                    BLF 410
C IERR   INTEGER                                                BLF 420
C IERR IS AN ERROR INDICATOR. ON EXIT                            BLF 430
C IERR = 0 IF NO ERROR HAS OCCURRED                             BLF 440
C IERR = 1 IF ON ENTRY IA.LT.N                                  BLF 450
C IERR = 2 IF THE MATRIX HAS A ROW CONSISTING ENTIRELY         BLF 460
C OF ZEROS                                                       BLF 470
C IERR = 3 IF THE MATRIX IS APPROXIMATELY SINGULAR              BLF 480
C IERR = 4 IF ON ENTRY N.LT.1                                   BLF 490
C                                                                    BLF 500
SUBROUTINE BLSOL(N, IR, TRANS, A, IA, P, B, IB, IERR)           BLS 10
C                                                                    BLS 20
C BLSOL CALCULATES THE APPROXIMATE SOLUTION OF A SET OF REAL    BLS 30
C LINEAR EQUATIONS WITH MULTIPLE RIGHT HAND SIDES A * X = B,    BLS 40
C AFTER A OR ITS TRANSPOSE HAS BEEN FACTORIZED BY BLCFAC.       BLS 50
C IT IS ORGANIZED SO THAT BLOCKS OF CONSECUTIVE RIGHT HAND     BLS 60
C SIDES ARE TREATED TOGETHER FOR EFFICIENCY ON A PAGED MACHINE. BLS 70
C                                                                    BLS 80
C PARAMETERS-                                                    BLS 90
C                                                                    BLS 100
C N      INTEGER                                                BLS 110
C ON ENTRY, N SPECIFIES THE ORDER OF THE MATRIX A                BLS 120
C (N.GT.0). UNCHANGED ON EXIT.                                  BLS 130
C                                                                    BLS 140
C IR     INTEGER                                                BLS 150
C ON ENTRY, IR SPECIFIES THE NUMBER OF RIGHT HAND SIDES        BLS 160
C I.E. THE NUMBER OF COLUMNS OF THE MATRIX B (IR.GT.0).      BLS 170

```

C	UNCHANGED ON EXIT.	BLS	180
C		BLS	190
C	TRANS LOGICAL	BLS	200
C	ON ENTRY, TRANS MUST BE TRUE IF THE TRANSPOSE OF A	BLS	210
C	WAS FACTORIZED BY BLC SOL AND FALSE OTHERWISE.	BLS	220
C	UNCHANGED ON EXIT.	BLS	230
C		BLS	240
C	A REAL ARRAY OF DIMENSION (IA,N)	BLS	250
C	BEFORE ENTRY, A MUST CONTAIN THE FACTORIZATION OF THE	BLS	260
C	MATRIX A OR ITS TRANSPOSE AS RETURNED BY THE	BLS	270
C	ROUTINE BLC SOL. UNCHANGED ON EXIT.	BLS	280
C		BLS	290
C	IA INTEGER	BLS	300
C	ON ENTRY, IA SPECIFIES THE FIRST DIMENSION OF A AS	BLS	310
C	DECLARED IN THE CALLING (SUB)PROGRAM (IA.GE.N).	BLS	320
C	UNCHANGED ON EXIT.	BLS	330
C		BLS	340
C	P REAL ARRAY OF DIMENSION (N)	BLS	350
C	BEFORE ENTRY, P MUST CONTAIN THE DETAILS OF THE ROW	BLS	360
C	PERMUTATIONS AS RETURNED BY THE ROUTINE BLCFAC	BLS	370
C		BLS	380
C	B REAL ARRAY OF DIMENSION (IB,IR)	BLS	390
C	BEFORE ENTRY, B MUST CONTAIN THE MATRIX OF RIGHT	BLS	400
C	HAND SIDES B.	BLS	410
C	ON SUCCESSFUL EXIT, B IS OVERWRITTEN BY THE MATRIX	BLS	420
C	OF SOLUTIONS X.	BLS	430
C		BLS	440
C	IB INTEGER	BLS	450
C	ON ENTRY, IB SPECIFIES THE FIRST DIMENSION OF B AS	BLS	460
C	DECLARED IN THE CALLING (SUB)PROGRAM (IB.GE.N).	BLS	470
C	UNCHANGED ON EXIT.	BLS	480
C		BLS	490
C	IERR INTEGER	BLS	500
C	IERR IS AN ERROR INDICATOR. ON EXIT	BLS	510
C	IERR = 0 IF NO ERROR HAS OCCURRED	BLS	520
C	IERR = 1 IF ON ENTRY IA.LT.N	BLS	530
C	IERR = 2 IF ON ENTRY IR.LT.1	BLS	540
C	IERR = 3 IF ON ENTRY IB.LT.N	BLS	550
C	IERR = 4 IF ON ENTRY N.LT.1	BLS	560
C		BLS	570

ALGORITHM 579

CPSC: Complex Power Series Coefficients

[D4]

BENGT FORNBERG
California Institute of Technology

Key Words and Phrases: numerical differentiation, Taylor series coefficients, analytic functions
CR Categories: 5.16
Language: FORTRAN

DESCRIPTION

An algorithm CPSC is presented here. It evaluates numerically the leading coefficients in a power series expansion of an analytic function (or, equivalently, a number of leading derivatives of an analytic function). A detailed description of the theoretical background of the code is given in [1], together with some warnings about cases in which full accuracy may not be reached.

Such cases are

- (1) very low-order polynomials (for example, $f(z) = 1 + z$);
- (2) functions whose Taylor expansions contain very large isolated terms (for example, $f(z) = 10^6 + (1/(1 - z))$);
- (3) certain entire functions (for example, $f(z) = e^z$);
- (4) functions whose radius of convergence is limited by a branch point at which the function remains many times differentiable (for example, $f(z) = (1 + z)^{10} \log(1 + z)$ expanded around $z = 0$).

In the case 1, the routine will normally fail to even approximate the correct answer. In cases 2, 3, and 4, problems are normally encountered only if large numbers of coefficients are wanted (e.g., more than 30). The supplied error estimate will, in most of these cases, give correct information about the lowered accuracy.

Inevitably there is a risk that the routine will attempt to evaluate the given function exactly at a singularity. This will happen, for example, for $f(z) = 1/(1 - z)$ expanded around zero if the initial radius is of the form $r = 2^p$, with p any integer. The risk for such exact coincidences in floating point can be minimized by choosing "irregular" initial radii. A practical procedure may be to start each case by the obtained final radius from a previous case.

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ALGORITHM

The user must enter the machine accuracy in a data statement. The routine adjusts the calculations accordingly in order to obtain the best possible accuracy. However, since a few decimal places are always lost in this routine, we do not recommend its general use if the machine accuracy is less than 10 significant decimal places. For example, for use on IBM System 370 machines, we recommend double or quadruple precision. The code below is for IBM double precision.

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```

C CPSC - COMPLEX POWER SERIES COEFFICIENTS
C
C BY BENGT FORNBERG
C
C ACM TRANSACTIONS ON MATHEMATICAL SOFTWARE, DECEMBER 1981
C
C     PROGRAM TEST (OUTPUT)
C
C THIS PROGRAM WITH THE FUNCTION F BELOW USES CPSC TO EVALUATE
C THE LEADING DERIVATIVES OF  $F(Z) = \text{CEXP}(Z) / (\text{CSIN}(Z)**3 + \text{CCOS}(Z)**3)$ 
C AT  $Z = \emptyset$ .
C     EXTERNAL F
C     DIMENSION A(51),IRANGE(4),ER(51)
C     COMPLEX A
C     DATA IRANGE/6,12,25,51/
C     DO 10 I=1,4
C         IR = IRANGE(I)
C         R = 1.
C         CALL CPSC(F, (0.,0.), IR,1,R,A,ER)
C         WRITE(6,20) I
C         WRITE(6,30) (A(J),J=1,IR)
10    WRITE(6,40) (ER(J),J=1,IR)
20    FORMAT(/20H DERIVATIVES , RANGE,I3)
30    FORMAT(4(1X,E18.10,E9.1))
40    FORMAT(/17H ESTIMATED ERRORS/(1X,16E8.1))
C     STOP
C     END
C     COMPLEX FUNCTION F(Z)
C
C TEST FUNCTION FOR USE WITH CPSC.
C
C     COMPLEX Z
C     F = CEXP(Z)/(CSIN(Z)**3+CCOS(Z)**3)
C     RETURN
C     END
C
C     SUBROUTINE CPSC(F,Z,N,IC,R,RS,ER)
C
C EVALUATION OF COMPLEX POWER SERIES COEFFICIENTS OR DERIVATIVES.
C
C *** INPUT PARAMETERS ***
C F COMPLEX FUNCTION, OF WHICH THE COEFFICIENTS OR DERIVATIVES
C ARE SOUGHT. THIS FUNCTION MUST BE DECLARED EXTERNAL IN THE
C CALLING PROGRAM.
C Z COMPLEX POINT AROUND WHICH F SHALL BE EXPANDED OR AT WHICH
C DERIVATIVES SHALL BE EVALUATED.
C N INTEGER, NUMBER OF COEFFICIENTS OR DERIVATIVES WANTED.
C N MUST BE GE 1 AND LE 51.
C IC SELECTS BETWEEN POWER SERIES COEFFICIENTS AND DREIVATIVES.
C IC .EQ. 0 ROUTINE RETURNS POWER SERIES COEFFICIENTS IN RS
C IC .NE. 0 ROUTINE RETURNS DERIVATIVES IN RS .
C *** INPUT AND OUTPUT PARAMETER ***
C R INITIAL RADIUS USED IN SEARCH FOR OPTIMAL RADIUS. THE RESULTING

```

```

C      RADIUS IS LEFT IN R. THE PROVIDED GUESS MAY BE IN ERROR WITH AT
C      MOST A FACTOR OF 3.E4 .
C      *** OUTPUT PARAMETERS ***
C      RS COMPLEX ARRAY RS(N) CONTAINING THE N FIRST
C      COEFFICIENTS (CORRESPONDING TO THE POWERS 0 TO N-1) OR DERIVA-
C      TIVES (ORDERS 0 TO N-1) .
C      ER REAL ARRAY ER(N) CONTAINING ERROR ESTIMATES FOR THE
C      NUMBERS IN RS(N).
C
C      DIMENSION IP(32),A(64),RS(N),ER(N),RT(51,3),FV(6),
C      * IW(7),SC(4),RV(3),C(4),FC(3)
C      COMPLEX F,A,V,RS,RT,FV,U,W,T,Z,RV,RQ,S,XK,MULT,CO
C
C      LIST OF THE VARIABLES INITIALIZED IN THE DATA STATEMENT BELOW.
C      EPS MACHINE ACCURACY. THIS CONSTANT HAS TO BE SUPPLIED BY
C      THE USER. IN THIS LIST, IT IS GIVEN AS 1.E-14 CORRESPONDING
C      TO THE 48 BIT FLOATING POINT MANTISSAS ON CDC CYBER MACHINES.
C      IND INTEGER FLAG.
C      L2 INTEGER FLAG.
C      IW 2**( 0 , 1 , 2 , 3 , 4 , 5 , 6 ) .
C      IP PERMUTATION CONSTANTS FOR THE FFT.
C      RV CONSTANTS FOR THE LAURENT SERIES TEST .
C
C      DATA EPS/1.E-14/,IND/0/,L2/1/,IW/1,2,4,8,16,32,64/,
C      * IP/64,32,48,16,56,24,40,8,60,28,44,12,52,20,36,4,62,30,46,14,
C      * 54,22,38,6,58,26,42,10,50,18,34,2/,
C      * RV/(-.4,.3),(.7,.2),(.02,-.06)/
C
C      STATEMENT FUNCTION FOR MULTIPLICATION OF A COMPLEX NUMBER
C      BY A REAL NUMBER.
C
C      MULT(RE,CO) = CMPLX(RE*REAL(CO),RE*AIMAG(CO))
C
C      EVALUATE SOME CONSTANTS THE FIRST TIME THE CODE IS EXECUTED.
C
C      IF(IND.EQ.1) GOTO 20
C      IND = 1
C      SC(1) = .125
C      C(1) = EPS**(1./28.)
C      EP6 = C(1)**6
C      PI = 4.*ATAN(1.)
C      FV(1) = (-1.,0.)
C      FV(2) = (0.,-1.)
C      R1 = SQRT(.5)
C      RA = 1./R1
C      FV(3) = CMPLX(R1,-R1)
C      DO 10 I=2,4
C          SC(I) = .5*SC(I-1)
C          C(I) = SQRT(C(I-1))
C          ANG = PI*SC(I-1)
C      10 FV(I+2) = CMPLX(COS(ANG),-SIN(ANG))
C      20 CONTINUE
C
C      START EXECUTION.
C
C      IF(N.GT.51.OR.N.LT.1) GOTO 260
C      LF = 0
C      NP = 0
C      M = 0
C      NR = -1
C
C      FIND IF A FFT OVER 8, 16, 32, OR 64 POINTS SHOULD BE USED.
C
C      KL = 1
C      IF(N.GT.6) KL=2
C      IF(N.GT.12) KL=3
C      IF(N.GT.25) KL=4
C      KM = KL+2
C      KN = 7-KM
C      IX = IW(KM+1)
C      IS = IW(KN)
C      30 V = CMPLX(R,0.)
C

```

```

C FUNCTION VALUES OF F ARE STORED READY PERMUTATED FOR THE FFT.
C
  DO 40 I=IS,32,IS
    IQ = IP(I)
    V = V*FV(KM)
    A(IQ) = F(Z+V)
  40 A(IQ-1) = F(Z-V)
    LN = 2
    JN = 1
C
C THE LOOP DO 70 ... CONSTITUTES THE FFT.
C
  DO 70 L=1,KM
    U = (1.,0.)
    W = FV(L)
    DO 60 J=1,JN
      DO 50 I=J,IX,LN
        IT = I+JN
        T = A(IT)*U
        A(IT) = A(I)-T
      50 A(I) = A(I)+T
      60 U = U*W
        LN = LN+LN
      70 JN = JN+JN
    CX = 0.
    B = 1.
C
C TEST ON HOW FAST THE COEFFICIENTS OBTAINED DECREASE.
C
  DO 80 I=1,IX
    CT = CABS(A(I))/B
    IF(CT.LT.CX) GOTO 80
    CX = CT
    INR = I
  80 B = B*C(KL)
    IF(M.LE.1) GOTO 100
C
C ESTIMATE OF THE ROUNDING ERROR LEVEL FOR THE LAST RADIUS.
C
  ER(1) = CX*EPS
  DO 90 I=2,N
  90 ER(I) = ER(I-1)/R
  100 SF = SC(KL)
  DO 110 I=1,IX
    A(I) = MULT(SF,A(I))
  110 SF = SF/R
    L1 = L2
    L2 = 1
    IF(INR.GT.IW(KL)) GOTO 150
    IF(LF.EQ.1) GOTO 140
C
C TEST IF THE SERIES IS A TAYLOR OR A LAURENT SERIES.
C
  SR = 0.
  SP = 0.
  DO 130 J=1,3
    RQ = MULT(R,RV(J))
    S = A(IX)
    DO 120 I=2,IX
      IA = IX+1-I
      S = S*RQ+A(IA)
  120 CP = CABS(S)
    IF(CP.GT.SP) SP=CP
    CM = CABS(S-F(Z+RQ))
  130 IF(CM.GT.SR) SR=CM
    IF(SR.GT.1.E-3*SP) GOTO 150
    LF = 1
  140 L2 = -1
C
C DETERMINATION OF THE NEXT RADIUS TO BE USED.
C
  150 IF(NR.GE.0) GOTO 160
    FACT = 2.

```

```

        IF(L2.EQ.1) FACT=.5
        L1 = L2
        NR = 0
160 IF(L1.NE.L2) GOTO 180
        IF(NR.GT.0) GOTO 170
        NP = NP+1
        IF(NP-15) 190,190,260
170 FACT = 1./FACT
180 FACT = 1./SQRT(FACT)
        NR = NR+1
190 R = R*FACT
        M = NR-KL-1
        IF(M.LE.0) GOTO 30
C
C THE RESULTS FOR THE LAST THREE RADII ARE STORED.
C
        DO 200 I=1,N
200 RT(I,M) = A(I)
        IF(M.EQ.1) GOTO 220
C
C EXTRAPOLATION.
C
        DO 210 I=1,N
            XK = RT(I,M-1)-RT(I,M)
210 RT(I,M-1) = RT(I,M)-MULT(FC(M-1),XK)
            IF(M.EQ.3) GOTO 230
C
C CALCULATION OF THE EXTRAPOLATION CONSTANTS.
C
220 FC(M) = 1.5+SIGN(.5,FACT-1.)
        IF(M.EQ.2) FC(M)=FC(M)+RA
        IF(FACT.GT.1.) FC(M)=-FC(M)
        GOTO 30
230 FC(3) = FC(1)*FC(2)/(FC(1)+FC(2)+1.)
C
C FINAL EXTRAPOLATION AND ERROR ESTIMATE.
C
        DO 240 I=1,N
            XK = RT(I,1)-RT(I,2)
            ER(I) = ER(I)+EP6*CABS(XK)
240 RS(I) = RT(I,2)-MULT(FC(3),XK)
C
C MULTIPLY POWER SERIES COEFFICIENTS AND ERROR ESTIMATE BY FACTORIALS
C IF DERIVATIVES WANTED.
C
        IF(IC.EQ.0) RETURN
        FAC = 0.
        FACT = 1.
        DO 250 I=1,N
            RS(I) = MULT(FACT,RS(I))
            ER(I) = FACT*ER(I)
            FAC = FAC+1.
250 FACT = FACT*FAC
        RETURN
C
C ERROR RETURN.
C
260 DO 270 I=1,N
            RS(I) = (0.,0.)
270 ER(I) = 1.E10
        RETURN
        END

```


ALGORITHM 580

QRUP: A Set of FORTRAN Routines for Updating QR Factorizations [F5]

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Key Words and Phrases: matrix factorization, orthogonalization
CR Categories: 5.14
Language: FORTRAN

DESCRIPTION

This algorithm is a translation into FORTRAN of the ALGOL routines published by Daniel et al. [2]. Given an $m \times n$ matrix A with $m \geq n$, these routines compute a QR factorization of A based on the Gram-Schmidt orthogonalization process, and they provide for updating the factorization when rows or columns are added to or deleted from A . Also the factors may be updated when A is modified by the addition or subtraction of a rank one matrix.

The translation is in most instances straightforward, but there are some points that need to be noted. Most of these are discussed in detail in the descriptive routine DESCRB, which is part of the algorithm.

However, a few comments belong here. Notice that DESCRB contains both a complete description of the routines and a fairly lengthy set of tests to ensure the accuracy of the translation. The tests basically consist of an initial factorization of a 7×7 Hilbert matrix, followed by random deletion of rows and columns of the matrix until it is just 1×1 , after which the rows and columns are, in a random order, reinserted to recreate the original matrix. Rank one matrices are introduced (then deleted) at several points as well. At each step, the factorization is updated and its accuracy is checked against the matrix being factored. A copy of the output from an execution of the tests follows the algorithm. The routines have also been tested through their use in an implementation of the minimization algorithm described in [1].

The translation is written in a portable subset of American National Standard FORTRAN, 1966, and has been thoroughly checked by the PFORT verifier. Except for some minor points noted in the routine DESCRB, the code is machine independent.

The routines are written simultaneously in single and double precision. As listed, compilation will give a single-precision version. It is straightforward,

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however, to obtain a double-precision version: just follow the simple instructions in DESCRB to interchange a few well-marked statements.

Notice that the basic linear algebra operations, such as taking norms or computing inner products, are done via calls to the BLAS described by Lawson et al. [3]. These routines are available through the International Mathematical and Statistical Libraries, Inc.,¹ or through SIAM.²

ACKNOWLEDGMENT

The author wishes to indicate his appreciation of the support received from NSERC, Canada.

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¹ NBC Building, 7500 Bellaire, Houston, TX 77036.

² 117 South 17 St., Philadelphia, PA 19103.

ALGORITHM

[A complete listing of the algorithm is available from the ACM Algorithms Distribution Service.]

ACM Transactions on Mathematical Software, Vol. 8, No. 4, December 1982, Page 405.

REMARK ON ALGORITHM 580

QRUP: A Set of FORTRAN Routines for Updating QR Factorizations [A. Buckley, *ACM Trans. Math. Softw.* 7, 4 (Dec. 1981), 548-549]

A. Buckley [Received 20 May 1982; accepted 20 May 1982]
Department of Management Information Systems, University of Arizona, Tucson, AZ 85721.

In two subprograms the functions **DNRM2()** and **DDOT()** are not declared to be **DOUBLE PRECISION** as they should be.

In program unit **DESCRB()**, sequence numbers **QRUP 139** and **QRUP 140** should be modified to read:

```
REAL          A, ASS, SDOT, COL, GLRMAX, OMEGA
C!!!! DOUBLE PRECISION  A, ASS, DDOT, COL, GLRMAX, OMEGA
```

In program unit **ORTCOL()**, sequence numbers **QRUP 762** and **QRUP 763** should be modified to read:

```
REAL          SDOT, SNRM2, OMEGA, ONE, ONENEG, Q, RHO
C!!!! DOUBLE PRECISION  DDOT, DNRM2, OMEGA, ONE, ONENEG, Q, RHO
```

ALGORITHM 581

An Improved Algorithm for Computing the Singular Value Decomposition [F1]

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Categories and Subject Descriptors: D.3.2 [Programming Languages]: Language Classifications—FORTRAN; G.1.3 [Numerical Analysis]: Numerical Linear Algebra—pseudoinverses; G.1.6 [Numerical Analysis]: Optimizations—least squares methods

General Terms: Algorithms

Additional Key Words and Phrases: singular value decomposition

DESCRIPTION

The set of FORTRAN subroutines given here is an implementation of the algorithm [1] for computing the Singular Value Decomposition (SVD) of a general m by n rectangular matrix A defined as

$$A = UWV^T,$$

where U is an $m \times \min(m, n)$ matrix containing the left singular vectors, W is a diagonal matrix of size $\min(m, n)$ containing the singular values, and V is an $n \times \min(m, n)$ matrix containing the right singular vectors. Note that m is allowed to be greater than or less than n . For ease of presentation, we assume m to be greater than or equal to n in the following discussion.

The algorithm is an improvement of the Golub-Reinsch algorithm [4], which is implemented in subroutines SVD and MINFIT in EISPACK [3] and in subroutine SSVDC in LINPACK [2]. It should be more efficient than the Golub-Reinsch algorithm when m is approximately larger than $2n$, as is the case in many least squares applications.

The algorithm has a hybrid nature. When m is about equal to n , the Golub-Reinsch algorithm is employed. When the ratio m/n is larger than a threshold value, which is determined by detailed operation counts [1], the improved algorithm is used.

The improved algorithm first computes the QR factorization of A using Householder transformations, and then uses the Golub-Reinsch algorithm on R . A further improvement over the Golub-Reinsch algorithm is when the left singular vectors are to be accumulated and saved. Here, instead of accumulating the Givens transformations (in the second phase of the algorithm where the singular values of the bidiagonal matrix are computed) on the $m \times n$ matrix containing the left singular vectors, we accumulate them on a temporary $n \times n$ matrix. This requires a small overhead in storage of an $n \times n$ matrix (small compared with $m \times n$) but offers big savings in time.

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An additional feature of the new algorithm is that it can accumulate all the left orthogonal transformations on a number of given vectors, which can then be used in computing least squares solutions. In this fashion, it is similar to the EISPACK routine MINFIT.

There are three main routines in the package:

- HYBSVD: This is the main routine which implements the hybrid algorithm.
 MGNSVD: This performs the same thing as HYBSVD, except that it assumes $m \geq n$.
 GRSVD: This is a slightly modified version of routine SVD in EISPACK which implements the Golub-Reinsch algorithm.

Besides, there are two utility routines:

- SSWAP: BLAS routine for swapping two vectors.
 SRELPR: Routine for computing the machine relative precision.

These five routines must be used together. They have been tested extensively on the IBM 370/168, 360/91 at the Stanford Linear Accelerator Center, and on the DEC 2060 in the Computer Science Department at Yale. They produce results that agree (up to machine precision) with those produced by SVD, MINFIT, and SSVDC. They have been verified by PFORT verifier [5] for portability.

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ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service (see page 91 for order form).]

```

SUBROUTINE HYBSVD(NA, NU, NV, NZ, NB, M, N, A, W, MATU, U, MATV, HYB  10
* V, Z, B, IRHS, IERR, RV1) HYB  20
INTEGER NA, NU, NV, NZ, M, N, IRHS, IERR, MIN0 HYB  30
REAL A(NA,1), W(1), U(NU,1), V(NV,1), Z(NZ,1), B(NB,IRHS), RV1(1) HYB  40
LOGICAL MATU, MATV HYB  50
C HYB  60
C THIS ROUTINE IS A MODIFICATION OF THE GOLUB-REINSCH PROCEDURE (1) HYB  70
C T HYB  80
C FOR COMPUTING THE SINGULAR VALUE DECOMPOSITION A = UWV OF A HYB  90
C REAL M BY N RECTANGULAR MATRIX. U IS M BY MIN(M,N) CONTAINING HYB 100
C THE LEFT SINGULAR VECTORS, W IS A MIN(M,N) BY MIN(M,N) DIAGONAL HYB 110
C MATRIX CONTAINING THE SINGULAR VALUES, AND V IS N BY MIN(M,N) HYB 120
C CONTAINING THE RIGHT SINGULAR VECTORS. HYB 130
C HYB 140
C THE ALGORITHM IMPLEMENTED IN THIS HYB 150
C ROUTINE HAS A HYBRID NATURE. WHEN M IS APPROXIMATELY EQUAL TO N, HYB 160
C THE GOLUB-REINSCH ALGORITHM IS USED, BUT WHEN EITHER OF THE RATIOS HYB 170
C M/N OR N/M IS GREATER THAN ABOUT 2, HYB 180
C A MODIFIED VERSION OF THE GOLUB-REINSCH HYB 190
C ALGORITHM IS USED. THIS MODIFIED ALGORITHM FIRST TRANSFORMS A HYB 200
C T HYB 210
C INTO UPPER TRIANGULAR FORM BY HOUSEHOLDER TRANSFORMATIONS L HYB 220
C AND THEN USES THE GOLUB-REINSCH ALGORITHM TO FIND THE SINGULAR HYB 230
C VALUE DECOMPOSITION OF THE RESULTING UPPER TRIANGULAR MATRIX R. HYB 240
C WHEN U IS NEEDED EXPLICITLY IN THE CASE M.GE.N (OR V IN THE CASE HYB 250
C M.LT.N), AN EXTRA ARRAY Z (OF SIZE AT LEAST HYB 260
C MIN(M,N)**2) IS NEEDED, BUT OTHERWISE Z IS NOT REFERENCED HYB 270
C AND NO EXTRA STORAGE IS REQUIRED. THIS HYBRID METHOD HYB 280

```

C	SHOULD BE MORE EFFICIENT THAN THE GOLUB-REINSCH ALGORITHM WHEN	HYB 290
C	M/N OR N/M IS LARGE. FOR DETAILS, SEE (2).	HYB 300
C		HYB 310
C	WHEN M .GE. N,	HYB 320
C	HYBSVD CAN ALSO BE USED TO COMPUTE THE MINIMAL LENGTH LEAST	HYB 330
C	SQUARES SOLUTION TO THE OVERDETERMINED LINEAR SYSTEM A*X=B.	HYB 340
C	IF M .LT. N (I.E. FOR UNDERDETERMINED SYSTEMS), THE RHS B	HYB 350
C	IS NOT PROCESSED.	HYB 360
C		HYB 370
C	NOTICE THAT THE SINGULAR VALUE DECOMPOSITION OF A MATRIX	HYB 380
C	IS UNIQUE ONLY UP TO THE SIGN OF THE CORRESPONDING COLUMNS	HYB 390
C	OF U AND V.	HYB 400
C		HYB 410
C	THIS ROUTINE HAS BEEN CHECKED BY THE PFORT VERIFIER (3) FOR	HYB 420
C	ADHERENCE TO A LARGE, CAREFULLY DEFINED, PORTABLE SUBSET OF	HYB 430
C	AMERICAN NATIONAL STANDARD FORTRAN CALLED PFORT.	HYB 440
C		HYB 450
C	REFERENCES:	HYB 460
C		HYB 470
C	(1) GOLUB,G.H. AND REINSCH,C. (1970) 'SINGULAR VALUE	HYB 480
C	DECOMPOSITION AND LEAST SQUARES SOLUTIONS,'	HYB 490
C	NUMER. MATH. 14,403-420, 1970.	HYB 500
C		HYB 510
C	(2) CHAN,T.F. (1982) 'AN IMPROVED ALGORITHM FOR COMPUTING	HYB 520
C	THE SINGULAR VALUE DECOMPOSITION,' ACM TOMS, VOL.8,	HYB 530
C	NO. 1, MARCH, 1982.	HYB 540
C		HYB 550
C	(3) RYDER,B.G. (1974) 'THE PFORT VERIFIER,' SOFTWARE -	HYB 560
C	PRACTICE AND EXPERIENCE, VOL.4, 359-377, 1974.	HYB 570
C		HYB 580
C	ON INPUT:	HYB 590
C		HYB 600
C	NA MUST BE SET TO THE ROW DIMENSION OF THE TWO-DIMENSIONAL	HYB 610
C	ARRAY PARAMETER A AS DECLARED IN THE CALLING PROGRAM	HYB 620
C	DIMENSION STATEMENT. NOTE THAT NA MUST BE AT LEAST	HYB 630
C	AS LARGE AS M.	HYB 640
C		HYB 650
C	NU MUST BE SET TO THE ROW DIMENSION OF THE TWO-DIMENSIONAL	HYB 660
C	ARRAY U AS DECLARED IN THE CALLING PROGRAM DIMENSION	HYB 670
C	STATEMENT. NU MUST BE AT LEAST AS LARGE AS M.	HYB 680
C		HYB 690
C	NV MUST BE SET TO THE ROW DIMENSION OF THE TWO-DIMENSIONAL	HYB 700
C	ARRAY PARAMETER V AS DECLARED IN THE CALLING PROGRAM	HYB 710
C	DIMENSION STATEMENT. NV MUST BE AT LEAST AS LARGE AS N.	HYB 720
C		HYB 730
C	NZ MUST BE SET TO THE ROW DIMENSION OF THE TWO-DIMENSIONAL	HYB 740
C	ARRAY PARAMETER Z AS DECLARED IN THE CALLING PROGRAM	HYB 750
C	DIMENSION STATEMENT. NOTE THAT NZ MUST BE AT LEAST	HYB 760
C	AS LARGE AS MIN(M,N).	HYB 770
C		HYB 780
C	NB MUST BE SET TO THE ROW DIMENSION OF THE TWO-DIMENSIONAL	HYB 790
C	ARRAY PARAMETER B AS DECLARED IN THE CALLING PROGRAM	HYB 800
C	DIMENSION STATEMENT. NB MUST BE AT LEAST AS LARGE AS M.	HYB 810
C		HYB 820
C	M IS THE NUMBER OF ROWS OF A (AND U).	HYB 830
C		HYB 840
C	N IS THE NUMBER OF COLUMNS OF A (AND NUMBER OF ROWS OF V).	HYB 850
C		HYB 860
C	A CONTAINS THE RECTANGULAR INPUT MATRIX TO BE DECOMPOSED.	HYB 870
C		HYB 880
C	B CONTAINS THE IRHS RIGHT-HAND-SIDES OF THE OVERDETERMINED	HYB 890
C	LINEAR SYSTEM A*X=B. IF IRHS .GT. 0 AND M .GE. N,	HYB 900
C	THEN ON OUTPUT, THE FIRST N COMPONENTS OF THESE IRHS COLUMNS	HYB 910
C	T	HYB 920
C	WILL CONTAIN U B. THUS, TO COMPUTE THE MINIMAL LENGTH LEAST	HYB 930
C	SQUARES SOLUTION, ONE MUST COMPUTE V*W TIMES THE COLUMNS OF	HYB 940
C	B, WHERE W IS A DIAGONAL MATRIX, W (I)=0 IF W(I) IS	HYB 950
C	NEGLIGIBLE, OTHERWISE IS 1/W(I). IF IRHS=0 OR M.LT.N,	HYB 960
C	B IS NOT REFERENCED.	HYB 970
C		HYB 980
C		HYB 990
C		HYB 1000
C	IRHS IS THE NUMBER OF RIGHT-HAND-SIDES OF THE OVERDETERMINED	HYB 1010
C	SYSTEM A*X=B. IRHS SHOULD BE SET TO ZERO IF ONLY THE SINGULAR	HYB 1020
C	VALUE DECOMPOSITION OF A IS DESIRED.	HYB 1030
C		HYB 1040

ALGORITHM 582

The Gibbs-Poole-Stockmeyer and Gibbs-King Algorithms for Reordering Sparse Matrices

JOHN G. LEWIS

Boeing Computer Services Co.

Categories and Subject Descriptors: G.1.3 [Numerical Analysis]: Numerical Linear Algebra—*sparse and very large systems*; G.m [Mathematics of Computing]: Miscellaneous—*FORTRAN*

General Terms: Algorithms

Additional Key Words and Phrases: Matrix bandwidth, matrix profile, matrix wavefront, banded matrix, Gibbs-Poole-Stockmeyer algorithm, Gibbs-King algorithm

DESCRIPTION

Given the structure of a symmetric or structurally symmetric sparse matrix, GPSKCA attempts to find a symmetric reordering of the matrix that produces a smaller bandwidth or profile. References [1], [4], [5], and [6] explain in detail the algorithms realized by GPSKCA. This algorithm provides the same mathematical capabilities as provided by REDUCE, Algorithms 508, and 509, but requires less memory and time and removes some implicit restrictions on the matrices that can be reordered. A description of the differences in the implementation and their effects is given in [7]; GPSKCA and REDUCE produce the same bandwidth and profile on all problems for which REDUCE executes successfully.

The package of subroutines is evoked by the FORTRAN statement

```
CALL GPSKCA (N, NZ, CONNEX, RSTART, DEGREE, OPTPRO, PERMUT, WORK,  
            WRKLEN, ERROR, SPACE)
```

where the parameters are described in the listing given here. The subroutines attempt to find a symmetric (row and column) permutation that reduces the bandwidth or profile of the reordered matrix. Whether to emphasize profile reduction or bandwidth reduction is determined by the logical parameter OPTPRO.

A common use for this algorithm will be prior to the solution of sparse linear algebraic equations or algebraic eigenvalue problems of moderately large size. The algorithm is supplied together with a FORTRAN program that accepts a sparse symmetric matrix in a standard sparse format and produces the equivalent reordered matrix in a format suitable for solving linear algebraic equations with the banded solvers in LINPACK [2] or the envelope solver in SPARSPAK [3].

The algorithm has been written in 1966 American National Standard FORTRAN. The package consists of 17 subroutines, GPSKCA, GPSKCB, . . . ,

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GPSKCQ. Special comment cards have been included so that the user can easily convert the code to an IBM FORTRAN version that uses half-length integers (INTEGER * 2) for all of the arrays except RSTART. This reduces the memory requirements by nearly a factor of 2 at the cost of restricting the order of the matrix to 32,767 or smaller. There would be no immediate restriction on the number of nonzeros in the matrix.

The primary difference in the usage of GPSKCA and its predecessor, REDUCE, is in the representation of the structure of the matrix. The format used herein is often much more efficient in space than that used in REDUCE. For the convenience of those users of REDUCE who would like to obtain the faster execution speeds of GPSKCA, but who cannot conveniently convert to the new, compact, format, an alternative version, GPSKRA, which accepts the older format for the matrix, will be supplied by the author upon request (not by the ACM distribution service).

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ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```

C === SEPARATOR === BEGINNING OF GPS AND GK ALGORITHMS
SUBROUTINE GPSKCA (N, DEGREE, RSTART, CONNED, OPTPRO, WRKLEN, GPSKCA 1
1 PERMUT, WORK, BANDWD, PROFIL, ERROR, SPACE)GPSKCA 2
C GPSKCA 3
C =====GPSKCA 4
C =====GPSKCA 5
C =GPSKCA 6
C = BANDWIDTH OR PROFILE REDUCTION =GPSKCA 7
C = FOR A SPARSE AND (STRUCTURALLY) SYMMETRIC MATRIX, =GPSKCA 8
C = USING EITHER =GPSKCA 9
C =GPSKCA10
C = THE GIBBS-POOLE-STOCKMEYER ALGORITHM (BANDWIDTH REDUCTION) =GPSKCA11
C = OR =GPSKCA12
C = THE GIBBS-KING ALGORITHM (PROFILE REDUCTION) =GPSKCA13
C =GPSKCA14
C =====GPSKCA15
C =====GPSKCA16
C = THIS CODE SUPERSEDES TOMS ALGORITHMS 508 AND 509 IN THE =GPSKCA17
C = COLLECTED ALGORITHMS OF THE ACM (CALGO). =GPSKCA18
C =====GPSKCA19
C =====GPSKCA20
C GPSKCA21
C -----GPSKCA22
C P A R A M E T E R S GPSKCA23
C -----GPSKCA24
C GPSKCA25
C INTEGER N, RSTART(N), WRKLEN, BANDWD, PROFIL, ERROR, SPACE GPSKCA26
C GPSKCA27
CIBM INTEGER *2 DEGREE(N), CONNED(1), PERMUT(N), WORK(WRKLEN), GPSKCA28
INTEGER DEGREE(N), CONNED(1), PERMUT(N), WORK(WRKLEN) GPSKCA29
C GPSKCA30
    
```



```

LOGICAL      OPTPRO
C
C -----
C
C INPUT PARAMETERS:
C -----
C
C      N      -- THE DIMENSION OF THE MATRIX
C
C      DEGREE,
C      RSTART,
C      CONNEC -- DESCRIBE THE STRUCTURE OF THE SPARSE MATRIX.
C                DEGREE(I) SPECIFIES THE NUMBER OF NON-ZERO
C                OFF-DIAGONAL ENTRIES IN THE I-TH ROW OF THE
C                SPARSE MATRIX. THE COLUMN INDICES OF THESE
C                ENTRIES ARE GIVEN IN CONSECUTIVE LOCATIONS IN
C                CONNEC, STARTING AT LOCATION RSTART(I).
C                IN OTHER WORDS, THE INDICES OF THE NON-ZERO
C                OFF-DIAGONAL ELEMENTS OF THE I-TH ROW ARE FOUND
C                IN:
C                CONNEC (RSTART(I)),
C                CONNEC (RSTART(I) + 1),
C                . . .
C                CONNEC (RSTART(I) + DEGREE(I) - 1)
C
C                DIMENSIONS:
C                RSTART IS DIMENSION N (OR LONGER).
C                DEGREE IS DIMENSION N (OR LONGER).
C                CONNEC IS DIMENSION ROUGHLY THE NUMBER OF NON-
C                ZERO ENTRIES IN THE MATRIX.
C
C      OPTPRO -- .TRUE. IF REDUCING THE PROFILE OF THE MATRIX
C                IS MORE IMPORTANT THAN REDUCING THE
C                BANDWIDTH
C                .FALSE. IF BANDWIDTH REDUCTION IS MOST IMPORTANT
C
C      WRKLEN -- THE ACTUAL LENGTH OF THE VECTOR WORK AS SUPPLIED
C                BY THE USER. SEE THE DISCUSSION OF THE WORKSPACE
C                'WORK' BELOW FOR TYPICAL STORAGE REQUIREMENTS.
C                THE VALUE OF WRKLEN WILL BE USED TO ENSURE THAT
C                THE ROUTINE WILL NOT USE MORE STORAGE THAN IS
C                AVAILABLE. IF NOT ENOUGH SPACE IS GIVEN IN WORK
C                TO PERMIT A SOLUTION TO BE FOUND, THE ERROR FLAG
C                WILL BE SET AND FURTHER COMPUTATION STOPPED.
C
C INPUT AND OUTPUT PARAMETER:
C -----
C
C      PERMUT -- ON INPUT, AN ALTERNATIVE REORDERING FOR THE
C                ROWS AND COLUMNS OF THE MATRIX. PERMUT(I) GIVES
C                THE POSITION IN WHICH ROW AND COLUMN I SHOULD
C                BE PLACED TO REDUCE THE BANDWIDTH OR THE PROFILE.
C                IF THE USER HAS NO ALTERNATIVE TO THE NATURAL
C                ORDERING IMPLICIT IN DEGREE, RSTART AND CONNEC,
C                HE SHOULD INITIALIZE PERMUT TO BE THE IDENTITY
C                PERMUTATION PERMUT(I) = I .
C
C                ON OUTPUT, PERMUT WILL CONTAIN THE PERMUTATION
C                FOR REORDERING THE ROWS AND COLUMNS WHICH REDUCES
C                THE BANDWIDTH AND/OR PROFILE. THE RESULT WILL BE
C                THE REORDERING FOUND BY 'GPSKCA' OR THE REORDERING
C                GIVEN BY THE USER IN 'PERMUT', WHICHEVER DOES THE
C                JOB BETTER.
C
C OUTPUT PARAMETERS:
C -----
C
C      WORK -- A TEMPORARY STORAGE VECTOR, OF LENGTH SOMEWHAT
C                GREATER THAN 3N. THE SPACE BEYOND 3N REQUIRED
C                IS PROBLEM-DEPENDENT. ANY PROBLEM CAN BE SOLVED
C                IN 6N+3 LOCATIONS.

```

GPSKCA31
GPSKCA32
GPSKCA33
GPSKCA34
GPSKCA35
GPSKCA36
GPSKCA37
GPSKCA38
GPSKCA39
GPSKCA40
GPSKCA41
GPSKCA42
GPSKCA43
GPSKCA44
GPSKCA45
GPSKCA46
GPSKCA47
GPSKCA48
GPSKCA49
GPSKCA50
GPSKCA51
GPSKCA52
GPSKCA53
GPSKCA54
GPSKCA55
GPSKCA56
GPSKCA57
GPSKCA58
GPSKCA59
GPSKCA60
GPSKCA61
GPSKCA62
GPSKCA63
GPSKCA64
GPSKCA65
GPSKCA66
GPSKCA67
GPSKCA68
GPSKCA69
GPSKCA70
GPSKCA71
GPSKCA72
GPSKCA73
GPSKCA74
GPSKCA75
GPSKCA76
GPSKCA77
GPSKCA78
GPSKCA79
GPSKCA80
GPSKCA81
GPSKCA82
GPSKCA83
GPSKCA84
GPSKCA85
GPSKCA86
GPSKCA87
GPSKCA88
GPSKCA89
GPSKCA90
GPSKCA91
GPSKCA92
GPSKCA93
GPSKCA94
GPSKCA95
GPSKCA96
GPSKCA97
GPSKCA98
GPSKCA99
GPSK100
GPSK101
GPSK102
GPSK103

C MOST PROBLEMS CAN BE REORDERED WITH 4N LOCATIONS GPSKC104
 C IN 'WORK'. IF SPACE IS NOT A CONSTRAINT, PROVIDE GPSKC105
 C 6N+3 LOCATIONS IN 'WORK'. OTHERWISE, PROVIDE AS GPSKC106
 C MUCH MORE THAN 3N AS IS CONVENIENT AND CHECK THE GPSKC107
 C ERROR FLAG AND SPACE REQUIRED PARAMETERS (SEE BELOW)GPSKC108
 C GPSKC109
 C ON OUTPUT, THE 1ST N LOCATIONS OF WORK WILL BE GPSKC110
 C A LISTING OF THE ORIGINAL ROW AND COLUMN INDICES AS GPSKC111
 C THEY APPEAR IN THE COMPUTED REORDERING. GPSKC112
 C LOCATIONS N+1, . . . , 2N OF WORK WILL CONTAIN GPSKC113
 C THE NEW POSITIONS FOR THE EQUATIONS IN THE ORDER GPSKC114
 C FOUND BY GPSKCA. THUS, THE TWO VECTORS ARE INVERSE GPSKC115
 C PERMUTATIONS OF EACH OTHER. IF THE ORDERING GPSKC116
 C FOUND BY THIS ALGORITHM IS BETTER THAN THE USER- GPSKC117
 C SUPPLIED ORDER, THE SECOND PERMUTATION VECTOR IS GPSKC118
 C IDENTICAL TO THE RESULT RETURNED IN 'PERMUT'. GPSKC119
 C GPSKC120
 C BANDWD -- THE BANDWIDTH OF THE MATRIX WHEN ROWS AND COLUMNS GPSKC121
 C ARE REORDERED IN THE ORDERING RETURNED IN PERMUT. GPSKC122
 C GPSKC123
 C PROFIL -- THE PROFILE OF THE MATRIX WHEN ROWS AND COLUMNS ARE GPSKC124
 C REORDERED IN THE ORDERING RETURNED IN PERMUT. GPSKC125
 C GPSKC126
 C ERROR -- WILL BE EQUAL TO ZERO IF A NEW NUMBERING COULD BE GPSKC127
 C FOUND IN THE SPACE PROVIDED. OTHERWISE, ERROR GPSKC128
 C WILL BE SET TO A POSITIVE ERROR CODE (SEE TABLE GPSKC129
 C GIVEN BELOW). IF THE REORDERING ALGORITHM HAS BEEN GPSKC130
 C STOPPED BY LACK OF WORKSPACE, THE SPACE PARAMETER GPSKC131
 C WILL BE SET TO THE NUMBER OF ADDITIONAL LOCATIONS GPSKC132
 C REQUIRED TO COMPLETE AT LEAST THE NEXT PHASE OF GPSKC133
 C THE ALGORITHM. GPSKC134
 C GPSKC135
 C WHENEVER A NON-ZERO VALUE FOR ERROR IS GIVEN GPSKC136
 C PERMUT WILL RETAIN THE VALUES PROVIDED BY THE USER GPSKC137
 C AND THE SCALARS BANDWD AND PROFIL WILL BE SET TOGPSKC138
 C OUTRAGEOUS VALUES. IT IS THE USER'S RESPONSIBILITY GPSKC139
 C TO CHECK THE STATUS OF ERROR. GPSKC140
 C GPSKC141
 C SPACE -- WILL INDICATE EITHER HOW MUCH SPACE THE REORDERING GPSKC142
 C ACTUALLY REQUIRED OR HOW MUCH SPACE WILL BE GPSKC143
 C REQUIRED TO COMPLETE THE NEXT PHASE OF THE GPSKC144
 C REORDERING ALGORITHM. THE POSSIBLE OUTCOMES ARE .. GPSKC145
 C GPSKC146
 C ERROR = 0 SPACE IS THE MINIMAL VALUE FORGPSKC147
 C WRKLEN REQUIRED TO REORDER GPSKC148
 C THIS MATRIX AGAIN. GPSKC149
 C GPSKC150
 C ERROR FJ 0 SPACE IS THE MINIMUM NUMBER GPSKC151
 C DUE TO LACK OF OF EXTRA WORKSPACE REQUIRED. GPSKC152
 C WORKSPACE TO CONTINUE THE REORDERING GPSKC153
 C ALGORITHM ON THIS MATRIX. GPSKC154
 C GPSKC155
 C ERROR FJ 0 SPACE = -1 GPSKC156
 C DUE TO ERROR GPSKC157
 C IN DATA STRUCTURES GPSKC158
 C GPSKC159
 C GPSKC160
 C -----GPSKC161
 C GPSKC162
 C -----GPSKC163
 C E R R O R C O D E S GPSKC164
 C -----GPSKC165
 C GPSKC166
 C ERROR CODES HAVE THE FORM 0XY OR 1XY. GPSKC167
 C GPSKC168
 C ERRORS OF THE FORM 1XY RESULT FROM INADEQUATE WORKSPACE. GPSKC169
 C GPSKC170
 C ERRORS OF THE FORM 0XY ARE INTERNAL PROGRAM CHECKS, WHICH GPSKC171
 C MOST LIKELY OCCUR BECAUSE THE CONNECTIVITY STRUCTURE OF THE GPSKC172
 C MATRIX IS REPRESENTED INCORRECTLY (E.G., THE DEGREE OF GPSKC173
 C A NODE IS NOT CORRECT OR NODE I IS CONNECTED TO NODE J, GPSKC174
 C BUT NOT CONVERSELY). GPSKC175

ALGORITHM 583

LSQR: Sparse Linear Equations and Least Squares Problems

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Categories and Subject Descriptors: G.1.3 [Numerical Analysis]: Numerical Linear Algebra—*linear systems (direct and iterative methods)*; G.3 [Mathematics of Computing]: Probability and Statistics—*statistical computing; statistical software*; G.m [Mathematics of Computing]: Miscellaneous—*FORTRAN program units*

General Terms: Algorithms

Additional Key Words and Phrases: Analysis of variance, conjugate-gradient method, least squares, linear equations, regression, sparse matrix

1. INTRODUCTION

LSQR finds a solution x to the following problems:

$$\text{Unsymmetric equations: solve } Ax = b \quad (1.1)$$

$$\text{Linear least squares: minimize } \|Ax - b\|_2 \quad (1.2)$$

$$\text{Damped least squares: minimize } \left\| \begin{bmatrix} A \\ \lambda I \end{bmatrix} x - \begin{bmatrix} b \\ 0 \end{bmatrix} \right\|_2 \quad (1.3)$$

where A is a matrix with m rows and n columns, b is an m -vector, λ is a scalar, and the given data A , b , λ are real. The matrix A will normally be large and sparse. It is defined by means of a user-written subroutine APROD, whose essential function is to compute products of the form Ax and $A^T y$ for given vectors x and y .

Problems (1.1) and (1.2) are treated as special cases of (1.3), which we shall write as

$$\min \| \bar{A} x - \bar{b} \|_2, \quad \bar{A} = \begin{bmatrix} A \\ \lambda I \end{bmatrix}, \quad \bar{b} = \begin{bmatrix} b \\ 0 \end{bmatrix}. \quad (1.4)$$

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Table I. Comparison of CGLS and LSQR

	Storage	Work per iteration
CGLS, $\lambda = 0$	$2m + 2n$	$2m + 3n$
CGLS, $\lambda \neq 0$	$2m + 2n$	$2m + 5n$
LSQR, any λ	$m + 2n$	$3m + 5n$

An earlier successful method for such problems is the *conjugate-gradient method for least squares systems* given by Hestenes and Stiefel [3]. (This method is described as algorithm CGLS in [6, sect. 7.1].) CGLS and LSQR are iterative methods with similar qualitative properties. Their computational requirements are summarized in Table I. In addition they require a product Ax and a product $A^T y$ each iteration.

In order to achieve the storage shown for LSQR, we ask the user to implement the matrix-vector products in the form

$$y \leftarrow y + Ax \quad \text{and} \quad x \leftarrow x + A^T y, \quad (1.5)$$

where \leftarrow means that one of the given vectors is overwritten by the expression shown. (A parameter specifies which expression the user's subroutine APROD should compute on any given entry.) We see that LSQR has a storage advantage if the operations (1.5) can be performed with no additional storage beyond that required to represent A . For least squares applications with many observations ($m \gg n$), this could be useful.

The work shown in Table I is the number of floating-point multiplications per iteration, excluding the work involved in the products Ax , $A^T y$. Since CGLS is somewhat more efficient, we would not discourage using that method whenever A or \bar{A} is well conditioned. However, LSQR is likely to obtain a more accurate solution in fewer iterations if \bar{A} is moderately or severely ill-conditioned.

Let $\bar{r}_k = \bar{b} - \bar{A}x_k$ be the residual vector associated with the k th iteration. LSQR provides estimates of $\|x_k\|_2$, $\|\bar{r}_k\|_2$, $\|\bar{A}^T \bar{r}_k\|_2$, the norm of \bar{A} , the condition number of \bar{A} , and standard errors for the components of x . The last two items require a further $2n$ multiplications per iteration and an additional n -vector of storage.

Subroutine LSQR is written in the PFORT subset of American National Standard FORTRAN. It contains no machine-dependent constants. Auxiliary routines required are APROD, NORMLZ, SCOPY, SNRM2, and SSCAL. The last three correspond to members of the BLAS collection [5].

2. MATHEMATICAL BACKGROUND

Algorithmic details are given in [6], mainly for the case $\lambda = 0$. We summarize these here with λ reintroduced, and show that a given value of λ may be dealt with at negligible cost. The vector norm $\|v\|_2 = (v^T v)^{1/2}$ is used throughout.

LSQR uses an algorithm of Golub and Kahan to reduce A to lower bidiagonal form. The quantities produced from A and b after $k + 1$ steps of the bidiagonalization (procedure Bidiag 1 [6]) are

$$\begin{aligned} \beta_1 &= \|b\|, \\ U_{k+1} &= [u_1, u_2, \dots, u_{k+1}], \\ V_{k+1} &= [v_1, v_2, \dots, v_{k+1}], \end{aligned} \quad B_k = \begin{bmatrix} \alpha_1 & & & & & & \\ \beta_2 & \alpha_2 & & & & & \\ & \beta_3 & \alpha_3 & & & & \\ & & \beta_4 & \alpha_4 & & & \\ & & & \beta_5 & \alpha_5 & & \\ & & & & \beta_6 & \alpha_6 & \\ & & & & & \beta_7 & \alpha_7 \\ & & & & & & \beta_8 & \alpha_8 \\ & & & & & & & \beta_9 & \alpha_9 \\ & & & & & & & & \beta_{10} & \alpha_{10} \end{bmatrix}. \quad (2.1)$$

The k th approximation to the solution x is then defined to be $x_k = V_k y_k$, where y_k solves the subproblem

$$\min \left\| \begin{bmatrix} B_k \\ \lambda I \end{bmatrix} y_k - \begin{bmatrix} \beta_1 e_1 \\ 0 \end{bmatrix} \right\|. \quad (2.2)$$

Letting the associated residual vectors be

$$\begin{aligned} t_{k+1} &= \beta_1 e_1 - B_k y_k \\ r_k &= b - A x_k \\ \bar{r}_k &= \bar{b} - \bar{A} x_k, \end{aligned} \quad (2.3)$$

we find that the relations

$$\begin{aligned} r_k &= U_{k+1} t_{k+1} \\ A^T r_k &= \lambda^2 x_k + \alpha_{k+1} \tau_{k+1} U_{k+1} \end{aligned} \quad (2.4)$$

will hold to machine accuracy, where τ_{k+1} is the last component of t_{k+1} , and we therefore conclude that (r_k, x_k) will be an acceptable solution of (1.4) if the computed value of either $\|t_{k+1}\|$ or $|\alpha_{k+1} \tau_{k+1}|$ is suitably small.

Björck [1] has previously observed that subproblem (2.2) is the appropriate generalization of $\min \|B_k y_k - \beta_1 e_1\|$, when $\lambda \neq 0$. He also discusses methods for computing y_k and x_k efficiently for various λ and k .

In LSQR we assume that a single value of λ is given, and to save storage and work, we do not compute y_k, r_k , or t_{k+1} . The orthogonal factorization

$$Q_k \begin{bmatrix} B_k & \beta_1 e_1 \\ \lambda I & 0 \end{bmatrix} = \begin{bmatrix} R_k & f_k \\ 0 & \bar{\phi}_{k+1} \\ 0 & q_k \end{bmatrix} \quad (2.5)$$

is computed ($Q_k^T Q_k = I$; R_k upper bidiagonal, $k \times k$) and this would give $R_k y_k = f_k$, but instead we solve $R_k^T D_k^T = V_k^T$ and form $x_k = D_k f_k$.

The factorization (2.5) is formed similarly to the case $\lambda = 0$ in [6], except that two rotations are required per step instead of one. For $k = 2$, the factorization proceeds according to

$$\begin{aligned} \begin{bmatrix} \alpha_1 & & \beta_1 \\ \beta_2 & \alpha_2 & \\ & \beta_3 & \\ \lambda & & \lambda \end{bmatrix} &\rightarrow \begin{bmatrix} \tilde{\rho}_1 & & \tilde{\phi}_1 \\ \beta_2 & \alpha_2 & \\ & \beta_3 & \\ & & \lambda & \psi_1 \end{bmatrix} \rightarrow \begin{bmatrix} \rho_1 & \theta_2 & \phi_1 \\ & \bar{\rho}_2 & \bar{\phi}_2 \\ & \beta_3 & \\ & & \lambda & \psi_1 \end{bmatrix} \\ &\rightarrow \begin{bmatrix} \rho_1 & \theta_2 & \phi_1 \\ & \tilde{\rho}_2 & \tilde{\phi}_2 \\ & \beta_3 & \\ & & \lambda & \psi_1 \\ & & & \lambda & \psi_2 \end{bmatrix} \rightarrow \begin{bmatrix} \rho_1 & \theta_2 & \phi_1 \\ & \rho_2 & \phi_2 \\ & & \bar{\phi}_3 \\ & & \lambda & \psi_1 \\ & & & \lambda & \psi_2 \end{bmatrix}. \end{aligned}$$

Note that the first λ is rotated into the diagonal element α_1 . This alters the right-hand side $\beta_1 e_1$ to produce ψ_1 , the first component of q_k . An alternative is to rotate λ into β_2 (and similarly for later λ), since this does not affect the right-hand side and it more closely simulates the algorithm that results when LSQR is applied to \bar{A} and \bar{b} directly. However, the rotations then have a greater effect on B_k , and in practice the first option has proved to give marginally more accurate results.

The estimates required to implement the stopping criteria are

$$\begin{aligned} \|\bar{r}_k\|^2 &= \|r_k\|^2 + \lambda^2 \|x_k\|^2 \approx \bar{\phi}_{k+1}^2 + \|q_k\|^2, \\ \|\bar{A}^T \bar{r}_k\| &= \|A^T r_k - \lambda^2 x_k\| \approx \left| \frac{\alpha_{k+1} \beta_{k+1} \phi_k}{\rho_k} \right|. \end{aligned}$$

This is a simple generalization of the case $\lambda = 0$. No additional storage is needed for q_k , since only its norm is required. In short, although the presence of λ complicates the algorithm description, it adds essentially nothing to the storage and work per iteration.

3. REGULARIZATION AND RELATED WORK

Introducing λ as in (1.3) is just one way of "regularizing" the solution x , in the sense that it can reduce the size of the computed solution and make its components less sensitive to changes in the data. LSQR is applicable when a value of λ is known a priori. The value is entered via the subroutine parameter DAMP. A second method for regularizing x is available through LSQR's parameter ACOND, which can cause iterations to terminate before $\|x_k\|$ becomes large. A similar approach has recently been described by Wold et al. [9], who give an illuminating interpretation of the bidiagonalization as a partial least squares procedure. Their description will also be useful to those who prefer the notation of multiple regression.

Methods for choosing λ , and other approaches to regularization, are given in [1, 2, 4, 8] and elsewhere. For a philosophical discussion, see [7].

4. CODING APROD

The best way to compute $y + Ax$ and $x + A^T y$ depends upon the origin of the matrix A . We shall illustrate a case that commonly arises, in which A is a sparse matrix whose nonzero coefficients are stored by *rows* in a simple list. Let A have M rows, N columns, and NZ nonzeros. Conceptually we need three arrays dimensioned as REAL RA(NZ) and INTEGER JA(NZ), NA(M), where

RA(L) is the Lth nonzero of A , counting across row 1, then across row 2, and so on;

JA(L) is the column in which the Lth nonzero of A lies;

NA(I) is the number of nonzero coefficients in the Ith row of A .

These quantities may be used in a straightforward way, as shown in Figure 1 (a FORTRAN implementation). We assume that they are made available to APROD through COMMON, and that the actual array dimensions are suitably large.

Blank or labeled COMMON will often be convenient for transmitting data to APROD. (Of course, some of the data could be local to APROD.) For greater generality, the parameter lists for LSQR and APROD include two workspace arrays IW, RW and their lengths LENIW, LENRW. LSQR does not use these parameters directly; it just passes them to APROD.

Figure 2 illustrates their use on the same example (sparse A stored by rows). An auxiliary subroutine APROD1 is needed to make the code readable. A similar scheme should be used to initialize the workspace parameters prior to calling LSQR.

Returning to the example itself, it may often be natural to store A by *columns* rather than rows, using analogous data structures. However, we note that in sparse least squares applications, A may have many more rows than columns ($M \gg N$). In such cases it is vital to store A by rows as shown, if the machine being used has a paged (virtual) memory. Random access is then restricted to arrays of length N rather than M , and page faults will therefore be kept to a minimum.

Note also that the arrays RA, JA, NA are adequate for computing both Ax and $A^T y$; we do not need to store A by rows *and* by columns.

Regardless of the application, it will be apparent when coding APROD for the two values of MODE that the matrix A is effectively being defined *twice*. Great care must be taken to avoid coding inconsistent expressions $y + A_1 x$ and $x + A_2^T y$, where either A_1 or A_2 is different from the desired A . (If $A_1 \neq A_2$, algorithm LSQR will not converge.) Parameters ANORM, ACOND, and CONLIM provide a safeguard for such an event.

5. PRECONDITIONING

It is well known that conjugate-gradient methods can be accelerated if a nonsingular matrix M is available to approximate A in some useful sense. When A is

square and nonsingular, the system $Ax = b$ is equivalent to both of the following systems:

$$(M^{-1}A)x = c \quad \text{where } Mc = b; \quad (5.1)$$

$$(AM^{-1})z = b \quad \text{where } Mx = z. \quad (5.2)$$

For least squares systems (undamped), only the analogue of (5.2) is applicable:

$$\min \|Ax - b\|_2 = \min \|(AM^{-1})z - b\|_2, \quad \text{where } Mx = z. \quad (5.3)$$

```

SUBROUTINE APROD( MODE,M,N,X,Y,
*                LENIW,LENRW,IW,RW )
C
C   INTEGER    MODE,M,N,LENIW,LENRW
C   INTEGER    IW(LENIW)
C   REAL      X(N),Y(M),RW(LENRW)
C
C   APROD PERFORMS THE FOLLOWING FUNCTIONS:
C
C   IF MODE = 1, SET Y = Y + A*X
C   IF MODE = 2, SET X = X + A(TRANSPPOSE)*Y
C
C   WHERE A IS A MATRIX STORED BY ROWS IN
C   THE ARRAYS RA, JA, NA. IN THIS EXAMPLE,
C   RA, JA, NA ARE STORED IN COMMON.
C
C   REAL      RA
C   INTEGER    JA,NA
C   COMMON    RA(9000),JA(9000),NA(1000)
C
C   INTEGER    I,J,L,L1,L2
C   REAL      SUM,YI,ZERO
C
C   ZERO = 0.0
C   L2 = 0
C   IF (MODE .NE.1) GO TO 400
C
C   -----
C   MODE = 1 -- SET Y = Y + A*X.
C   -----
C   DO 200 I = 1, M
C     SUM = ZERO
C     L1 = L2 + 1
C     L2 = L2 + NA(I)
C     DO 100 L = L1, L2
C       J = JA(L)
C       SUM = SUM + RA(L)*X(J)
C   100 CONTINUE
C     Y(I) = Y(I) + SUM
C   200 CONTINUE
C     RETURN
C
C   -----
C   MODE = 2 -- SET X = X + A(TRANSPPOSE)*Y.
C   -----
C   400 DO 600 I = 1, M
C     YI = Y(I)
C     L1 = L2 + 1
C     L2 = L2 + NA(I)
C     DO 500 L = L1, L2
C       J = JA(L)
C       X(J) = X(J) + RA(L)*YI
C   500 CONTINUE
C   600 CONTINUE
C     RETURN
C
C   END OF APROD
C   END

```

Fig. 1. Computation of $y + Ax$, $x + A^T y$, where A is a sparse matrix stored compactly by rows. For convenience, the data structure for A is held in COMMON.

```

      SUBROUTINE APROD( MODE,M,N,X,Y,
*          LENIW,LENRW,IW,RW )
C
      INTEGER    MODE,M,N,LENIW,LENRW
      INTEGER    IW(LENIW)
      REAL      X(N),Y(M),RW(LENRW)
C
C   APROD PERFORMS THE FOLLOWING FUNCTIONS:
C
      IF MODE = 1, SET  Y = Y + A*X
      IF MODE = 2, SET  X = X + A(TRANSPPOSE)*Y
C
      WHERE A IS A MATRIX STORED BY ROWS IN
      THE ARRAYS RA, JA, NA. IN THIS EXAMPLE,
      APROD IS AN INTERFACE BETWEEN LSQR AND
      ANOTHER USER ROUTINE THAT DOES THE WORK.
      THE WORKSPACE ARRAY RW CONTAINS RA.
      THE FIRST M COMPONENTS OF IW CONTAIN NA,
      AND THE REMAINDER OF IW CONTAINS JA.
      THE DIMENSIONS OF RW AND IW ARE ASSUMED
      TO BE SUFFICIENTLY LARGE.
C
      INTEGER    LENJA,LENRA,LOCJA
C
      LOCJA = M + 1
      LENJA = LENIW - LOCJA + 1
      LENRA = LENRW
      CALL APROD1( MODE,M,N,X,Y,
*          LENJA,LENRA,IW,IW(LOCJA),RW )
      RETURN
C
      END OF APROD
      END

```

Fig. 2. Same as Figure 1, with the data structure for A held in the workspace parameters.

```

      SUBROUTINE APROD1( MODE,M,N,X,Y,
*          LENJA,LENRA,NA,JA,RA )
C
      INTEGER    MODE,M,N,LENJA,LENRA
      INTEGER    NA(M),JA(LENJA)
      REAL      X(N),Y(M),RA(LENRA)
C
C   APROD1 DOES THE WORK FOR APROD.
C
      INTEGER    I,J,L,L1,L2
      REAL      SUM,YI,ZERO
C
      < the same code as in APROD in Figure 1 >
C
      END OF APROD1
      END

```

```

LEAST-SQUARES TEST PROBLEM      P( 20 10 1 1 1.00E-03 )
CONDITION NO. = 9.9995E 00      RESIDUAL FUNCTION = 9.812157000E-01

```

```

LSQR  --  LEAST-SQUARES SOLUTION OF A*X = B
THE MATRIX A HAS 20 ROWS AND 10 COLS
THE DAMPING PARAMETER IS  DAMP = 1.00E-03
ATOL  = 1.00E-06          CONLIM = 1.00E 02
BTOL  = 1.00E-06          ITNLIM = 80

```

Fig. 3. Example output from test program and LSQR on a damped least squares problem.

ITN	X(1)	FUNCTION	COMPATIBLE	INCOMPATIBLE	NORM(ABAR)	COND(ABAR)
0	0.0000000000E-01	6.3410580000E 00	1.000E 00	9.135E-02		
1	-6.3564250000E-01	4.0387670000E 00	6.369E-01	7.244E-01	7.51E-01	1.00E 00
2	-4.7282630000E-01	2.3303970000E 00	3.675E-01	3.349E-01	1.10E 00	2.42E 00
3	-2.8075080000E-01	1.8962160000E 00	2.990E-01	2.462E-01	1.30E 00	3.63E 00
4	2.6825070000E-01	1.5905200000E 00	2.508E-01	1.424E-01	1.49E 00	5.27E 00
5	1.2649560000E 00	1.4032960000E 00	2.213E-01	1.160E-01	1.59E 00	7.09E 00
6	2.0648040000E 00	1.2910880000E 00	2.036E-01	9.273E-02	1.70E 00	9.03E 00
7	3.0031450000E 00	1.2072930000E 00	1.904E-01	8.294E-02	1.79E 00	1.12E 01
8	3.7526340000E 00	1.1551220000E 00	1.822E-01	5.138E-02	1.90E 00	1.34E 01
9	5.4443550000E 00	1.0780640000E 00	1.700E-01	3.155E-02	1.95E 00	1.72E 01
10	8.9918140000E 00	9.8151740000E-01	1.548E-01	1.283E-02	1.96E 00	2.44E 01
11	8.9998990000E 00	9.8120520000E-01	1.547E-01	2.017E-04	2.20E 00	2.75E 01
12	8.9999290000E 00	9.8120540000E-01	1.547E-01	4.361E-06	2.38E 00	2.98E 01
13	8.9999280000E 00	9.8120550000E-01	1.547E-01	9.078E-07	2.48E 00	3.13E 01

NO. OF ITERATIONS = 13 STOPPING CONDITION = 2

THE LEAST-SQRS SOLN IS GOOD ENOUGH, GIVEN ATOL

	RESIDUAL NORM (ABAR*X - BBAR)	RESIDUAL NORM (NORMAL EQNS)	SOLUTION NORM (X)						
ESTIMATED BY LSQR	9.812055E-01	2.206693E-06	1.688187E 01						
COMPUTED FROM X	9.812157E-01	1.083419E-05	1.688184E 01						
SOLUTION									
1	8.99993	2	7.99997	3	6.99998	4	5.99999	5	4.99999
6	3.99999	7	3.00000	8	2.00000	9	0.999994	10	-0.204206E-05
STANDARD ERRORS									
1	2.11589	2	0.888101	3	0.685644	4	0.556184	5	0.614104
6	0.409182	7	0.565480	8	0.519385	9	0.375466	10	0.589787

Figure 3 (continued)

We note only that subroutine *LSQR* may be applied without change to systems (5.1)-(5.3). The effect of M is localized to the user's own subroutine *APROD*. For example, when $\text{MODE} = 1$, *APROD* for the last two systems should compute $y + (AM^{-1})x$ by first solving $Mw = x$ and then computing $y + Aw$. Clearly it must be possible to solve systems involving M and M^T very efficiently.

6. OUTPUT

Subroutine *LSQR* produces printed output on file *NOUT*, if the parameter *NOUT* is positive. This is illustrated in Figure 3, in which the least squares problem solved is $P(20, 10, 1, 1)$ as defined in [6], with a slight generalization to include a damping parameter $\lambda = 10^{-3}$. (Single precision was used on an IBM 370/168.) The items printed at the k th iteration are as follows.

ITN	The iteration number k . Results are always printed for the first 10 and last 10 iterations. Intermediate results are printed if $m \leq 40$ or $n \leq 40$, or if one of the convergence conditions is nearly satisfied. Otherwise, information is printed every 10th iteration.
X(1)	The value of the first element of the approximate solution x_k .
FUNCTION	The value of the function being minimized, namely $\ \bar{r}_k\ = (\ r_k\ ^2 + \lambda^2 \ x_k\ ^2)^{1/2}$.
COMPATIBLE	A dimensionless quantity which should converge to zero if and only if $Ax = b$ is compatible. It is an estimate of $\ \bar{r}_k\ / \ b\ $, which decreases monotonically.

- INCOMPATIBLE A dimensionless quantity which should converge to zero *if and only if* the optimum $\|\bar{r}_k\|$ is nonzero. It is an estimate of $\|\bar{A}^T \bar{r}_k\| / (\|\bar{A}\|_F \|\bar{r}_k\|)$, which is usually *not* monotonic.
- NORM(ABAR) A monotonically increasing estimate of $\|\bar{A}\|_F$.
- COND(ABAR) A monotonically increasing estimate of $\text{cond}(\bar{A}) = \|\bar{A}\|_F \|\bar{A}^+\|_F$, the condition number of \bar{A} .

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ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```

SUBROUTINE LSQR( M,N,APROD,DAMP,
1          LENIW,LENRW,IW,RW,
2          U,V,W,X,SE,
3          ATOL,BTOL,CONLIM,ITNLIM,NOUT,
4          ISTOP,ANORM,ACOND,RNORM,ARNORM,XNORM )
C
EXTERNAL  APROD
INTEGER  M,N,LENIW,LENRW,ITNLIM,NOUT,ISTOP
INTEGER  IW(LENIW)
REAL    RW(LENRW),U(M),V(N),W(N),X(N),SE(N),
1      ATOL,BTOL,CONLIM,DAMP,ANORM,ACOND,RNORM,ARNORM,XNORM
C
-----
C
C
C  LSQR FINDS A SOLUTION X TO THE FOLLOWING PROBLEMS...
C
C  1. UNSYMMETRIC EQUATIONS -- SOLVE A*X = B
C
C  2. LINEAR LEAST SQUARES -- SOLVE A*X = B
C                               IN THE LEAST-SQUARES SENSE
C
C  3. DAMPED LEAST SQUARES -- SOLVE ( A ) * X = ( B )
C                               ( DAMP*I ) ( 0 )
C                               IN THE LEAST-SQUARES SENSE
C
C  WHERE A IS A MATRIX WITH M ROWS AND N COLUMNS, B IS AN
C  M-VECTOR, AND DAMP IS A SCALAR (ALL QUANTITIES REAL).
C  THE MATRIX A IS INTENDED TO BE LARGE AND SPARSE. IT IS ACCESSED
C  BY MEANS OF SUBROUTINE CALLS OF THE FORM

```

C
 C CALL APROD(MODE,M,N,X,Y,LENIW,LENRW,IW,RW) 29.
 C
 C WHICH MUST PERFORM THE FOLLOWING FUNCTIONS... 30.
 C 31.
 C 32.
 C 33.
 C IF MODE = 1, COMPUTE $Y = Y + A*X$. 34.
 C IF MODE = 2, COMPUTE $X = X + A(\text{TRANSPOSE})*Y$. 35.
 C 36.
 C THE VECTORS X AND Y ARE INPUT PARAMETERS IN BOTH CASES. 37.
 C IF MODE = 1, Y SHOULD BE ALTERED WITHOUT CHANGING X. 38.
 C IF MODE = 2, X SHOULD BE ALTERED WITHOUT CHANGING Y. 39.
 C THE PARAMETERS LENIW, LENRW, IW, RW MAY BE USED FOR WORKSPACE 40.
 C AS DESCRIBED BELOW. 41.
 C 42.
 C THE RHS VECTOR B IS INPUT VIA U, AND SUBSEQUENTLY OVERWRITTEN. 43.
 C 44.
 C 45.
 C NOTE. LSQR USES AN ITERATIVE METHOD TO APPROXIMATE THE SOLUTION. 46.
 C THE NUMBER OF ITERATIONS REQUIRED TO REACH A CERTAIN ACCURACY 47.
 C DEPENDS STRONGLY ON THE SCALING OF THE PROBLEM. POOR SCALING OF 48.
 C THE ROWS OR COLUMNS OF A SHOULD THEREFORE BE AVOIDED WHERE 49.
 C POSSIBLE. 50.
 C 51.
 C FOR EXAMPLE, IN PROBLEM 1 THE SOLUTION IS UNALTERED BY 52.
 C ROW-SCALING. IF A ROW OF A IS VERY SMALL OR LARGE COMPARED TO 53.
 C THE OTHER ROWS OF A, THE CORRESPONDING ROW OF (A B) SHOULD 54.
 C BE SCALED UP OR DOWN. 55.
 C 56.
 C IN PROBLEMS 1 AND 2, THE SOLUTION X IS EASILY RECOVERED 57.
 C FOLLOWING COLUMN-SCALING. IN THE ABSENCE OF BETTER INFORMATION, 58.
 C THE NONZERO COLUMNS OF A SHOULD BE SCALED SO THAT THEY ALL HAVE 59.
 C THE SAME EUCLIDEAN NORM (E.G. 1.0). 60.
 C 61.
 C IN PROBLEM 3, THERE IS NO FREEDOM TO RE-SCALE IF DAMP IS 62.
 C NONZERO. HOWEVER, THE VALUE OF DAMP SHOULD BE ASSIGNED ONLY 63.
 C AFTER ATTENTION HAS BEEN PAID TO THE SCALING OF A. 64.
 C 65.
 C THE PARAMETER DAMP IS INTENDED TO HELP REGULARIZE 66.
 C ILL-CONDITIONED SYSTEMS, BY PREVENTING THE TRUE SOLUTION FROM 67.
 C BEING VERY LARGE. ANOTHER AID TO REGULARIZATION IS PROVIDED BY 68.
 C THE PARAMETER ACOND, WHICH MAY BE USED TO TERMINATE ITERATIONS 69.
 C BEFORE THE COMPUTED SOLUTION BECOMES VERY LARGE. 70.
 C 71.
 C 72.
 C NOTATION 73.
 C ----- 74.
 C 75.
 C THE FOLLOWING QUANTITIES ARE USED IN DISCUSSING THE SUBROUTINE 76.
 C PARAMETERS... 77.
 C 78.
 C ABAR = (A), BBAR = (B) 79.
 C (DAMP*I) (0) 80.
 C 81.
 C R = B - A*X, RBAR = BBAR - ABAR*X 82.
 C 83.
 C RNORM = $\text{SQRT}(\text{NORM}(R)**2 + \text{DAMP}**2 * \text{NORM}(X)**2)$ 84.
 C = $\text{NORM}(\text{RBAR})$ 85.
 C 86.
 C RELPR = THE RELATIVE PRECISION OF FLOATING-POINT ARITHMETIC 87.
 C ON THE MACHINE BEING USED. FOR EXAMPLE, ON THE IBM 370, 88.
 C RELPR IS ABOUT $1.0E-6$ AND $1.0D-16$ IN SINGLE AND DOUBLE 89.
 C PRECISION RESPECTIVELY. 90.
 C 91.
 C LSQR MINIMIZES THE FUNCTION RNORM WITH RESPECT TO X. 92.
 C 93.
 C 94.
 C PARAMETERS 95.
 C ----- 96.
 C 97.
 C M INPUT THE NUMBER OF ROWS IN A. 98.
 C 99.
 C N INPUT THE NUMBER OF COLUMNS IN A. 100.

C				101.
C	APROD	EXTERNAL	SEE ABOVE.	102.
C				103.
C	DAMP	INPUT	THE DAMPING PARAMETER FOR PROBLEM 3 ABOVE.	104.
C			(DAMP SHOULD BE 0.0 FOR PROBLEMS 1 AND 2.)	105.
C			IF THE SYSTEM $A*X = B$ IS INCOMPATIBLE, VALUES	106.
C			OF DAMP IN THE RANGE 0 TO $\sqrt{\text{RELPR}}*\text{NORM}(A)$	107.
C			WILL PROBABLY HAVE A NEGLIGIBLE EFFECT.	108.
C			LARGER VALUES OF DAMP WILL TEND TO DECREASE	109.
C			THE NORM OF X AND TO REDUCE THE NUMBER OF	110.
C			ITERATIONS REQUIRED BY LSQR.	111.
C				112.
C			THE WORK PER ITERATION AND THE STORAGE NEEDED	113.
C			BY LSQR ARE THE SAME FOR ALL VALUES OF DAMP.	114.
C				115.
C	LENIW	INPUT	THE LENGTH OF THE WORKSPACE ARRAY IW.	116.
C	LENRW	INPUT	THE LENGTH OF THE WORKSPACE ARRAY RW.	117.
C	IW	WORKSPACE	AN INTEGER ARRAY OF LENGTH LENIW.	118.
C	RW	WORKSPACE	A REAL ARRAY OF LENGTH LENRW.	119.
C				120.
C			NOTE. LSQR DOES NOT EXPLICITLY USE THE PREVIOUS FOUR	121.
C			PARAMETERS, BUT PASSES THEM TO SUBROUTINE APROD FOR	122.
C			POSSIBLE USE AS WORKSPACE. IF APROD DOES NOT NEED	123.
C			IW OR RW, THE VALUES $\text{LENIW} = 1$ OR $\text{LENRW} = 1$ SHOULD	124.
C			BE USED, AND THE ACTUAL PARAMETERS CORRESPONDING TO	125.
C			IW OR RW MAY BE ANY CONVENIENT ARRAY OF SUITABLE TYPE.	126.
C				127.
C	U(M)	INPUT	THE RHS VECTOR B. BEWARE THAT U IS	128.
C			OVER-WRITTEN BY LSQR.	129.
C				130.
C	V(N)	WORKSPACE		131.
C	W(N)	WORKSPACE		132.
C				133.
C	X(N)	OUTPUT	RETURNS THE COMPUTED SOLUTION X.	134.
C				135.
C	SE(N)	OUTPUT	RETURNS STANDARD ERROR ESTIMATES FOR THE	136.
C			COMPONENTS OF X. FOR EACH I, SE(I) IS SET	137.
C			TO THE VALUE $\text{RNORM} * \sqrt{\text{SIGMA}(I,I) / T}$,	138.
C			WHERE $\text{SIGMA}(I,I)$ IS AN ESTIMATE OF THE I-TH	139.
C			DIAGONAL OF THE INVERSE OF $\text{ABAR}(\text{TRANPOSE})*\text{ABAR}$	140.
C			AND $T = 1$ IF $M \leq N$,	141.
C			$T = M - N$ IF $M > N$ AND $\text{DAMP} = 0$,	142.
C			$T = M$ IF $\text{DAMP} \neq 0$.	143.
C				144.
C	ATOL	INPUT	AN ESTIMATE OF THE RELATIVE ERROR IN THE DATA	145.
C			DEFINING THE MATRIX A. FOR EXAMPLE,	146.
C			IF A IS ACCURATE TO ABOUT 6 DIGITS, SET	147.
C			$\text{ATOL} = 1.0\text{E-}6$.	148.
C				149.
C	BTOL	INPUT	AN ESTIMATE OF THE RELATIVE ERROR IN THE DATA	150.
C			DEFINING THE RHS VECTOR B. FOR EXAMPLE,	151.
C			IF B IS ACCURATE TO ABOUT 6 DIGITS, SET	152.
C			$\text{BTOL} = 1.0\text{E-}6$.	153.
C				154.
C	CONLIM	INPUT	AN UPPER LIMIT ON $\text{COND}(\text{ABAR})$, THE APPARENT	155.
C			CONDITION NUMBER OF THE MATRIX ABAR.	156.
C			ITERATIONS WILL BE TERMINATED IF A COMPUTED	157.
C			ESTIMATE OF $\text{COND}(\text{ABAR})$ EXCEEDS CONLIM.	158.
C			THIS IS INTENDED TO PREVENT CERTAIN SMALL OR	159.
C			ZERO SINGULAR VALUES OF A OR ABAR FROM	160.
C			COMING INTO EFFECT AND CAUSING UNWANTED GROWTH	161.
C			IN THE COMPUTED SOLUTION.	162.
C				163.
C			CONLIM AND DAMP MAY BE USED SEPARATELY OR	164.
C			TOGETHER TO REGULARIZE ILL-CONDITIONED SYSTEMS.	165.
C				166.
C			NORMALLY, CONLIM SHOULD BE IN THE RANGE	167.
C			1000 TO $1/\text{RELPR}$.	168.
C			SUGGESTED VALUE --	169.
C			$\text{CONLIM} = 1/(100*\text{RELPR})$ FOR COMPATIBLE SYSTEMS,	170.
C			$\text{CONLIM} = 1/(10*\sqrt{\text{RELPR}})$ FOR LEAST SQUARES.	171.
C				172.

C			NOTE. IF THE USER IS NOT CONCERNED ABOUT THE PARAMETERS	173.
C			ATOL, BTOL AND CONLIM, ANY OR ALL OF THEM MAY BE SET	174.
C			TO ZERO. THE EFFECT WILL BE THE SAME AS THE VALUES	175.
C			RELPR, RELPR AND 1/RELPR RESPECTIVELY.	176.
C				177.
C	ITNLIM	INPUT	AN UPPER LIMIT ON THE NUMBER OF ITERATIONS.	178.
C			SUGGESTED VALUE --	179.
C			ITNLIM = N/2 FOR WELL CONDITIONED SYSTEMS,	180.
C			ITNLIM = 4*N OTHERWISE.	181.
C				182.
C	NOUT	INPUT	FILE NUMBER FOR PRINTER. IF POSITIVE,	183.
C			A SUMMARY WILL BE PRINTED ON FILE NOUT.	184.
C				185.
C	ISTOP	OUTPUT	AN INTEGER GIVING THE REASON FOR TERMINATION...	186.
C				187.
C		0	X = 0 IS THE EXACT SOLUTION.	188.
C			NO ITERATIONS WERE PERFORMED.	189.
C				190.
C		1	THE EQUATIONS A*X = B ARE PROBABLY	191.
C			COMPATIBLE. NORM(A*X - B) IS SUFFICIENTLY	192.
C			SMALL, GIVEN THE VALUES OF ATOL AND BTOL.	193.
C				194.
C		2	THE SYSTEM A*X = B IS PROBABLY NOT	195.
C			COMPATIBLE. A LEAST-SQUARES SOLUTION HAS	196.
C			BEEN OBTAINED WHICH IS SUFFICIENTLY ACCURATE,	197.
C			GIVEN THE VALUE OF ATOL.	198.
C				199.
C		3	AN ESTIMATE OF COND(ABAR) HAS EXCEEDED	200.
C			CONLIM. THE SYSTEM A*X = B APPEARS TO BE	201.
C			ILL-CONDITIONED. OTHERWISE, THERE COULD BE AN	202.
C			AN ERROR IN SUBROUTINE APROD.	203.
C				204.
C		4	THE EQUATIONS A*X = B ARE PROBABLY	205.
C			COMPATIBLE. NORM(A*X - B) IS AS SMALL AS	206.
C			SEEMS REASONABLE ON THIS MACHINE.	207.
C				208.
C		5	THE SYSTEM A*X = B IS PROBABLY NOT	209.
C			COMPATIBLE. A LEAST-SQUARES SOLUTION HAS	210.
C			BEEN OBTAINED WHICH IS AS ACCURATE AS SEEMS	211.
C			REASONABLE ON THIS MACHINE.	212.
C				213.
C		6	COND(ABAR) SEEMS TO BE SO LARGE THAT THERE IS	214.
C			NOT MUCH POINT IN DOING FURTHER ITERATIONS,	215.
C			GIVEN THE PRECISION OF THIS MACHINE.	216.
C			THERE COULD BE AN ERROR IN SUBROUTINE APROD.	217.
C				218.
C		7	THE ITERATION LIMIT ITNLIM WAS REACHED.	219.
C				220.
C	ANORM	OUTPUT	AN ESTIMATE OF THE FROBENIUS NORM OF ABAR.	221.
C			THIS IS THE SQUARE-ROOT OF THE SUM OF SQUARES	222.
C			OF THE ELEMENTS OF ABAR.	223.
C			IF DAMP IS SMALL AND IF THE COLUMNS OF A	224.
C			HAVE ALL BEEN SCALED TO HAVE LENGTH 1.0,	225.
C			ANORM SHOULD INCREASE TO ROUGHLY SQRT(N).	226.
C			A RADICALLY DIFFERENT VALUE FOR ANORM MAY	227.
C			INDICATE AN ERROR IN SUBROUTINE APROD (THERE	228.
C			MAY BE AN INCONSISTENCY BETWEEN MODES 1 AND 2).	229.
C				230.
C	ACOND	OUTPUT	AN ESTIMATE OF COND(ABAR), THE CONDITION	231.
C			NUMBER OF ABAR. A VERY HIGH VALUE OF ACOND	232.
C			MAY AGAIN INDICATE AN ERROR IN APROD.	233.
C				234.
C	RNORM	OUTPUT	AN ESTIMATE OF THE FINAL VALUE OF NORM(RBAR),	235.
C			THE FUNCTION BEING MINIMIZED (SEE NOTATION	236.
C			ABOVE). THIS WILL BE SMALL IF A*X = B HAS	237.
C			A SOLUTION.	238.
C				239.
C	ARNORM	OUTPUT	AN ESTIMATE OF THE FINAL VALUE OF	240.
C			NORM(ABAR(TRANSPOSE)*RBAR), THE NORM OF	241.
C			THE RESIDUAL FOR THE USUAL NORMAL EQUATIONS.	242.
C			THIS SHOULD BE SMALL IN ALL CASES. (ARNORM	243.
C			WILL OFTEN BE SMALLER THAN THE TRUE VALUE	244.

C		COMPUTED FROM THE OUTPUT VECTOR X.)	245.
C			246.
C	XNORM	OUTPUT AN ESTIMATE OF THE NORM OF THE FINAL	247.
C		SOLUTION VECTOR X.	248.
C			249.
C			250.
C	SUBROUTINES AND FUNCTIONS USED		251.
C	-----		252.
C			253.
C	USER	APROD	254.
C	LSQR	NORMLZ	255.
C	BLAS	SCOPY, SNRM2, SSCAL (SEE LAWSON ET AL. BELOW)	256.
C		(SNRM2 IS USED ONLY IN NORMLZ)	257.
C	FORTRAN	ABS, MOD, SQRT	258.
C			259.
C			260.
C	PRECISION		261.
C	-----		262.
C			263.
C		THE NUMBER OF ITERATIONS REQUIRED BY LSQR WILL USUALLY DECREASE	264.
C		IF THE COMPUTATION IS PERFORMED IN HIGHER PRECISION. TO CONVERT	265.
C		LSQR AND NORMLZ BETWEEN SINGLE- AND DOUBLE-PRECISION, CHANGE	266.
C		THE WORDS	267.
C		SCOPY, SNRM2, SSCAL	268.
C		ABS, REAL, SQRT	269.
C		TO THE APPROPRIATE BLAS AND FORTRAN EQUIVALENTS.	270.
C			271.
C			272.
C	REFERENCES		273.
C	-----		274.
C			275.
C	PAIGE, C.C. AND SAUNDERS, M.A.	LSQR: AN ALGORITHM FOR SPARSE	276.
C		LINEAR EQUATIONS AND SPARSE LEAST SQUARES.	277.
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C			279.
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C			285.
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C	-----		287.

ALGORITHM 584

CUBTRI: Automatic Cubature over a Triangle

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Categories and Subject Descriptors: G.1.4 [Numerical Analysis]: Quadrature and Numerical Differentiation—*adaptive quadrature, multiple quadrature*; G.m [Mathematics of Computing]: Miscellaneous—*FORTRAN*

General Terms: Algorithms, Theory

Additional Key Words and Phrases: Quadrature rule

DESCRIPTION AND PURPOSE

Given a triangle T with vertices (t_{1j}, t_{2j}) , $j = 1, 2, 3$, a function $f(x, y)$ defined on T , and a tolerance $\epsilon > 0$, CUBTRI attempts to compute a number A such that

$$\left| A - \int \int_T f(x, y) dx dy \right| < \max\{\epsilon, |A\epsilon|\}. \quad (1)$$

The method employed is similar to that of Haegemans [1] and consists of the following steps.

- (1) Compute an approximate integral A and error estimate η for the original triangle T .
- (2) If the current values of A and η satisfy $\eta < \max\{\epsilon, |A\epsilon|\}$, or if further computation will exceed specified limits, exit.
- (3) Given a list of triangles whose union is T , together with an approximate integral and error estimate for each triangle, identify the triangle T_k with largest error estimate.
- (4) Remove T_k from the data list and subtract its contributions from the current values of A and η .
- (5) Divide T_k into four congruent triangles, append them to the data list, and compute an approximate integral and error estimate for each. Add these contributions to A and η .
- (6) Diagnose whether round-off errors are contaminating the obtained values to a point where further subdivision is futile. If so, exit; else return to step 2.

The general strategy falls within the framework outlined by Rice [7], apart from the test for roundoff, which is due to Piessens [5] and also appears in Haegemans' routine TRIADA [1].

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Table I. The Integration Formula Pair Used by CUBTRI

i	$\zeta_1^{(i)}$	$\zeta_2^{(i)}$	$\zeta_3^{(i)}$	$w_1^{(i)}$	$w_2^{(i)}$
0	1/3	1/3	1/3	9/40	7137/62720 -45 σ /1568
1, 2	3/7 $\pm 2\phi/21$	2/7 $\mp \phi/21$	$=\zeta_2^{(i)}$	31/80 $\mp \phi/400$	-9301697/4695040 $\mp 13517313\phi/23475200$ +764885 σ /939008 $\pm 198763\phi\sigma/939008$
3, 4	4/9 $\pm \phi/9$ + $\sigma/9$ $\mp \sigma\phi/45$	5/18 $\mp \phi/18$ - $\sigma/18$ $\pm \sigma\phi/90$	$=\zeta_2^{(i)}$	0	102791225/59157504 $\pm 23876225\phi/59157504$ -34500875 σ /59157504 $\mp 9914825\phi\sigma/59157504$
5	$=\zeta_2^{(3)}$	$\zeta_2^{(4)}$	4/9 + $\sigma/9$	0	11075/8064 -125 σ /288

$$\phi = \sqrt{15}, \sigma = \sqrt{7}$$

$$Q^{(i)}f = \frac{1}{6} \sum_{k=1}^3 \sum_{l=1}^3 f(\zeta_k^{(i)}, \zeta_l^{(i)})$$

$$Q_1f = \Delta \sum_{i=0}^2 w_1^{(i)} Q^{(i)}f$$

$$Q_2f = \Delta \sum_{i=0}^5 w_2^{(i)} Q^{(i)}f$$

Δ = area of triangle

$\zeta_k^{(i)}$ are barycentric coordinates

CUBTRI attempts to improve on TRIADA in terms of number of function evaluations and storage management.

NUMBER OF FUNCTION EVALUATIONS

Two cubature rules Q_1 and Q_2 are applied to each triangle to produce estimates A_1 and A_2 , where A_2 is of a higher degree than A_1 . This gives an approximate integral $A = A_2$ and error estimate $\eta = |A_2 - A_1|$ which is likely to be pessimistic, since it really estimates the error in A_1 rather than that in A_2 . This effect is countered by employing a device given in [3] for sharpening the error estimate. It is, of course, possible to use two completely independent cubature rules (e.g., Haegemans [1] uses conic product Gauss formulas (see Stroud [8], Section 2.5) with 36 and 49 points, respectively), but it seems sensible to extend Kronrod's [2] idea, that Q_2 should reuse the points of Q_1 , to the two-dimensional case.

In CUBTRI, Q_1 is the 7-point degree-5 rule of Radon [6], and Q_2 is a new 19-point degree-8 rule given in Table I. The theory of Lyness and Jespersen [4] shows that 19 points and degree 8 are optimal, given that Q_2 has D_3 symmetry and contains the 7 points of Q_1 . A curious feature of Q_2 is that the 12 new points are the intersection points of three pairs of parallel lines (see Figure 1).

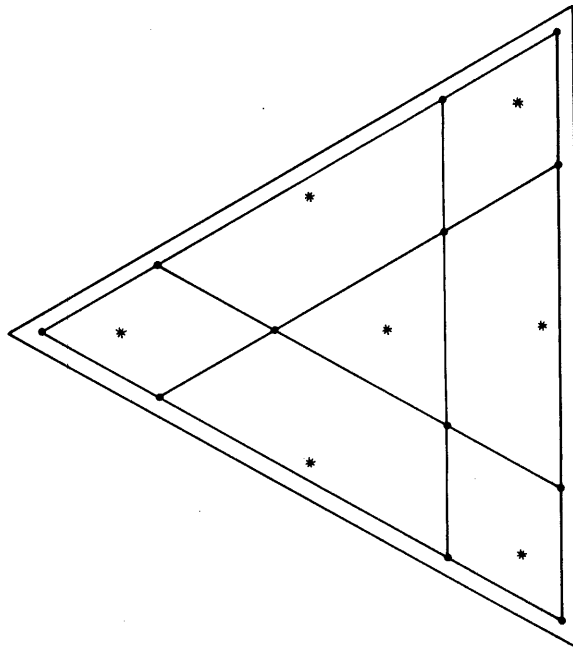
STORAGE MANAGEMENT

Most of the storage requirements arise from the need to store a list of triangles. In general, 6 numbers are necessary to describe a triangle (e.g., three vertices of two coordinates each), giving a total of $6n$ numbers for n triangles.

Haegemans [1] has improved on this by keeping a separate list of vertices; each triangle can then be described by 3 integers. The length of the list is only $n + 2$, because each subdivision creates as many new vertices (three) as the net increase in the number of triangles. This yields a total of $3n$ integers and $2n + 4$ reals.

In CUBTRI we take advantage of the fact that the sides of each subtriangle are parallel to those of the original triangle. Geometrically, one can visualize the subtriangle as being obtained by translation, scaling, and possibly rotation

Fig. 1. Disposition of points in the degree-8 formula. •—Additional points of degree-8 formula. *—Points of Radon's degree-5 formula.



through 180° . These effects can be represented by two coordinates, a scale factor, and the sign of the scale factor, respectively: thus three numbers per triangle suffice. It is convenient to use barycentric coordinates for the translation vector, because in that case only numbers that are terminating binary fractions are generated; these are exact machine numbers (even in single precision) on most computers, and rounding errors are avoided. Only $3n$ reals are therefore required in addition to the vertices of T .

TESTS

The driver main program and function $f(x, y)$ supplied with the algorithm calculate the following seven integrals.

Case 1.

$$\int_0^1 \int_0^x (x^2 + 3y^2)^{-1/2} dy dx = \frac{\log(2 + \sqrt{3})}{\sqrt{3}}.$$

This integrand has a singularity at the origin.

Case 2. The same as Case 1, with a rapidly oscillatory relative error of magnitude 10^{-6} perturbing the integrand. This simulates the behavior of a machine with a short word length.

Case 3.

$$\int_0^1 \int_0^x f(x, y) dy dx = \frac{\pi}{8}$$

where

$$f(x, y) = \begin{cases} 1, & (x - \frac{1}{2})^2 + (y - \frac{1}{2})^2 \leq \frac{1}{2}, \\ 0, & \text{otherwise.} \end{cases}$$

This integrand has a jump discontinuity along a curved line, the circle of radius $\frac{1}{2}$ centered at $(\frac{1}{2}, \frac{1}{2})$.

The following four test functions were suggested by Haegemans [1].

Case 4.

$$\int_0^1 \int_0^{1-x} \exp(\sin x \cos y) dy dx \approx 0.6918104506612316.$$

Table II. $f(x, y) = (x^2 + 3y^2)^{-1/2}$

ϵ	CUBTRI				TRIADA			
	NFE	EST	ERR	IER	NFE	EST	ERR	IER
$1_{10^{-3}}$	323	$6.2_{10^{-4}}$	$6.1_{10^{-4}}$	0	765	$9.1_{10^{-4}}$	$3.0_{10^{-3}}$	0
$1_{10^{-4}}$	551	$8.5_{10^{-5}}$	$7.5_{10^{-5}}$	0	2125	$5.7_{10^{-5}}$	$1.9_{10^{-4}}$	0
$1_{10^{-5}}$	1007	$8.1_{10^{-6}}$	$1.5_{10^{-6}}$	0	3145	$7.1_{10^{-6}}$	$2.3_{10^{-5}}$	0
$1_{10^{-6}}$	1919	$8.6_{10^{-7}}$	$3.7_{10^{-8}}$	0	4165	$9.0_{10^{-7}}$	$2.9_{10^{-6}}$	0
$1_{10^{-7}}$	3135	$9.9_{10^{-8}}$	$3.0_{10^{-9}}$	0	5525	$6.3_{10^{-8}}$	$1.8_{10^{-7}}$	0
$1_{10^{-8}}$	5035	$9.5_{10^{-9}}$	$1.2_{10^{-10}}$	0	7225	$9.5_{10^{-9}}$	$6.6_{10^{-9}}$	0
$1_{10^{-9}}$	7391	$9.8_{10^{-10}}$	$7.6_{10^{-11}}$	0	10285	$9.8_{10^{-10}}$	$9.8_{10^{-10}}$	0
						10		

Table III. $f(x, y) = (x^2 + 3y^2)^{-1/2}(1 + \eta)$, $|\eta| \leq 10^{-6}$

ϵ	CUBTRI				TRIADA			
	NFE	EST	ERR	IER	NFE	EST	ERR	IER
$1_{10^{-3}}$	323	$6.2_{10^{-4}}$	$6.1_{10^{-4}}$	0	765	$9.1_{10^{-4}}$	$3.0_{10^{-3}}$	0
$1_{10^{-4}}$	551	$8.5_{10^{-5}}$	$7.5_{10^{-5}}$	0	2125	$5.7_{10^{-5}}$	$1.9_{10^{-4}}$	0
$1_{10^{-5}}$	1007	$8.2_{10^{-6}}$	$1.5_{10^{-6}}$	0	3145	$7.2_{10^{-6}}$	$2.3_{10^{-5}}$	0
$1_{10^{-6}}$	1919	$8.6_{10^{-7}}$	$2.8_{10^{-8}}$	0	4165	$9.7_{10^{-7}}$	$2.9_{10^{-6}}$	0
$1_{10^{-7}}$	3135	$9.8_{10^{-8}}$	$1.1_{10^{-8}}$	0	6205	$1.0_{10^{-7}}$	$8.9_{10^{-9}}$	0
$1_{10^{-8}}$	5415	$1.3_{10^{-8}}$	$1.5_{10^{-8}}$	4	9945	$6.5_{10^{-8}}$	$1.7_{10^{-10}}$	2

Table IV. $f(x, y) = \frac{1}{2} + \frac{1}{2} \text{sign}\left\{\frac{1}{4} - \left(x - \frac{1}{2}\right)^2 - \left(y - \frac{1}{2}\right)^2\right\}$

ϵ	CUBTRI				TRIADA			
	NFE	EST	ERR	IER	NFE	EST	ERR	IER
$1_{10^{-3}}$	551	$3.9_{10^{-3}}$	$2.6_{10^{-3}}$	4	10625	$9.9_{10^{-4}}$	$6.4_{10^{-5}}$	0
$1_{10^{-4}}$	3059	$2.0_{10^{-3}}$	$5.7_{10^{-4}}$	4	14705	$9.3_{10^{-4}}$	$3.0_{10^{-4}}$	1

This integral seems to be impossible to evaluate analytically; the numerical "exact" value is given in [1].

Case 5.

$$\int_0^1 \int_0^{1-x} \frac{x^2}{(1+x^2)} dy dx = \frac{1 + \log 2}{2} - \frac{\pi}{4}.$$

The integrand depends on x only.

Case 6.

$$\int_0^1 \int_0^{1-x} \sin(3x + 6y) dy dx = \frac{1}{9} \sin 3 - \frac{1}{18} \sin 6.$$

A mildly oscillating integrand.

Case 7.

$$\int_0^1 \int_0^{1-x} \sin(x+y)^{-1/2} dy dx = \frac{2}{3}.$$

The integrand has a gentler singularity at the origin than that of Case 1.

In Tables II-VIII we report the results obtained by CUBTRI as well as by TRIADA [2] on a CDC Cyber 174 (48-bit mantissa, binary floating point). The column legends are

- ϵ Required tolerance.
- NFE Number of function evaluations.
- EST Estimated absolute error.
- ERR Actual absolute error.

Table V. $f(x, y) = \exp(\sin x \cos y)$

ϵ	CUBTRI				TRIADA			
	NFE	EST	ERR	IER	NFE	EST	ERR	IER
$1_{10^{-3}}$	19	$1.8_{10^{-5}}$	$2.6_{10^{-9}}$	0	85	$1.5_{10^{-10}}$	$2.3_{10^{-12}}$	0
$1_{10^{-5}}$	95	$4.0_{10^{-11}}$	$4.6_{10^{-12}}$	0	85	$1.5_{10^{-10}}$	$2.3_{10^{-12}}$	0

Table VI. $f(x, y) = x^2/(1 + x^2)$

ϵ	CUBTRI				TRIADA			
	NFE	EST	ERR	IER	NFE	EST	ERR	IER
$1_{10^{-3}}$	19	$3.2_{10^{-5}}$	$1.7_{10^{-7}}$	0	85	$8.6_{10^{-10}}$	$6.0_{10^{-11}}$	0
$1_{10^{-5}}$	95	$7.2_{10^{-10}}$	$3.6_{10^{-10}}$	0	85	$8.6_{10^{-10}}$	$6.0_{10^{-11}}$	0

Table VII. $f(x, y) = \sin(3x + 6y)$

ϵ	CUBTRI				TRIADA			
	NFE	EST	ERR	IER	NFE	EST	ERR	IER
$1_{10^{-3}}$	19	$3.3_{10^{-4}}$	$2.1_{10^{-8}}$	0	85	$1.9_{10^{-9}}$	$2.1_{10^{-11}}$	0
$1_{10^{-4}}$	95	$1.6_{10^{-9}}$	$1.3_{10^{-11}}$	0	85	$1.9_{10^{-9}}$	$2.1_{10^{-11}}$	0
$1_{10^{-9}}$	171	$8.7_{10^{-10}}$	$2.9_{10^{-11}}$	0	425	$2.4_{10^{-12}}$	$4.8_{10^{-15}}$	0

Table VIII. $f(x + y) = (x + y)^{-1/2}$

ϵ	CUBTRI				TRIADA			
	NFE	EST	ERR	IER	NFE	EST	ERR	IER
$1_{10^{-3}}$	95	$1.4_{10^{-4}}$	$1.4_{10^{-4}}$	0	85	$3.2_{10^{-4}}$	$6.2_{10^{-4}}$	0
$1_{10^{-4}}$	171	$4.9_{10^{-5}}$	$4.9_{10^{-5}}$	0	765	$4.8_{10^{-5}}$	$9.8_{10^{-5}}$	0
$1_{10^{-5}}$	323	$6.3_{10^{-6}}$	$6.1_{10^{-6}}$	0	1445	$6.0_{10^{-6}}$	$1.2_{10^{-5}}$	0
$1_{10^{-6}}$	475	$9.4_{10^{-7}}$	$7.7_{10^{-7}}$	0	2125	$7.6_{10^{-7}}$	$1.5_{10^{-6}}$	0
$1_{10^{-7}}$	1007	$7.2_{10^{-8}}$	$1.2_{10^{-8}}$	0	2805	$9.5_{10^{-8}}$	$1.9_{10^{-7}}$	0
$1_{10^{-8}}$	1615	$9.1_{10^{-9}}$	$1.5_{10^{-9}}$	0	3825	$4.2_{10^{-9}}$	$8.4_{10^{-9}}$	0
$1_{10^{-9}}$	2375	$8.8_{10^{-10}}$	$5.6_{10^{-11}}$	0	4505	$5.5_{10^{-10}}$	$1.1_{10^{-9}}$	0

IER Error flag. For TRIADA, IER = 1 denotes that the limit on NFE has been reached, and IER = 2 that the round-off error threshold has been reached. For CUBTRI, the meaning of IER is explained in the leading comments.

We can draw the following conclusions.

(1) *Reliability.* The error estimates produced by CUBTRI after subdivision range from 0.87 to 120 times the actual error, and in only one instance is the actual error underestimated. Before subdivision (i.e., when NFE = 19), the device [3] is unavailable, and the overestimation is much more pessimistic. The error estimates produced by TRIADA range from 0.3 to 500 times the actual error, and in 16 instances (all involving nonsmooth integrands) the actual error is underestimated.

(2) *Accuracy.* On smooth integrands TRIADA is slightly more accurate than CUBTRI (e.g., in Case 2, TRIADA with NFE = 85 produces ERR = $2.3_{10^{-12}}$). On nonsmooth integrands CUBTRI is much more accurate than TRIADA (e.g., in Case 1, CUBTRI with NFE = 1919 produces ERR = $3.7_{10^{-8}}$, whereas TRIADA with NFE = 2125 can only obtain ERR = $1.9_{10^{-4}}$).

(3) *Efficiency.* On smooth integrands TRIADA in 15 instances out of 21 terminated at a given ϵ with fewer function evaluations than CUBTRI. On nonsmooth integrands, in 20 instances out of 21 CUBTRI terminated with fewer function evaluations than TRIADA.

On machines with a short word length, termination with IER = 4 can be expected when ϵ is smaller than machine accuracy, and the behavior of the algorithm when ϵ is close to machine accuracy might be somewhat different. The situation is simulated by the difference between Cases 1 and 2.

To summarize, CUBTRI is a reliable and efficient integrator on smooth as well as nonsmooth integrands, although not quite so accurate and efficient for smooth integrands as TRIADA.

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The referee of an early version of this paper drew my attention to [4]. Elsie de Lange wrote the driver main program and test functions. Ina Marais checked the correctness of Table I.

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ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Distribution Service.]

```

SUBROUTINE CUBTRI(F, T, EPS, MCALLS, ANS, ERR, NCALLS, W, NW,      CUB  10
* IDATA, RDATA, IER)                                           CUB  20
C                                                                    CUB  30
C      ADAPTIVE CUBATURE OVER A TRIANGLE                          CUB  40
C                                                                    CUB  50
C      PARAMETERS                                                 CUB  60
C      F      - USER SUPPLIED EXTERNAL FUNCTION OF THE FORM     CUB  70
C              F(X,Y, IDATA, RDATA)                               CUB  80
C              WHERE X AND Y ARE THE CARTESIAN COORDINATES OF A   CUB  90
C              POINT IN THE PLANE, AND IDATA AND RDATA ARE INTEGER CUB 100
C              AND REAL VECTORS IN WHICH DATA MAY BE PASSED.    CUB 110
C      T      - ARRAY OF DIMENSION (2,3) WHERE T(1,J) AND T(2,J) CUB 120
C              ARE THE X AND Y COORDINATES OF THE J-TH VERTEX OF CUB 130
C              THE GIVEN TRIANGLE (INPUT)                          CUB 140
C      EPS    - REQUIRED TOLERANCE (INPUT). IF THE COMPUTED        CUB 150
C              INTEGRAL IS BETWEEN-1 AND 1, AN ABSOLUTE ERROR    CUB 160
C              TEST IS USED, ELSE A RELATIVE ERROR TEST IS USED. CUB 170
C      MCALLS- MAXIMUM PERMITTED NUMBER OF CALLS TO F (INPUT)    CUB 180
C      ANS    - ESTIMATE FOR THE VALUE OF THE INTEGRAL OF F OVER  CUB 190
C              THE GIVEN TRIANGLE (OUTPUT)                          CUB 200
C      ERR    - ESTIMATED ABSOLUTE ERROR IN ANS (OUTPUT)         CUB 210
C      NCALLS- ACTUAL NUMBER OF CALLS TO F (OUTPUT). THIS        CUB 220
C              PARAMETER MUST BE INITIALIZED TO 0 ON THE FIRST    CUB 230
C              CALL TO CUBTRI FOR A GIVEN INTEGRAL (INPUT)        CUB 240
C      W      - WORK SPACE. MAY NOT BE DESTROYED BETWEEN CALLS TO CUB 250
C              CUBTRI IF RESTARTING IS INTENDED                    CUB 260
C      NW     - LENGTH OF WORK SPACE (INPUT).                      CUB 270
C              IF NW .GE. 3*(19+3*MCALLS)/38, TERMINATION DUE TO  CUB 280
C              FULL WORK SPACE WILL NOT OCCUR.                     CUB 290
C      IER    - TERMINATION INDICATOR (OUTPUT)                     CUB 300

```

C	IER=0	NORMAL TERMINATION, TOLERANCE SATISFIED	CUB	310
C	IER=1	MAXIMUM NUMBER OF CALLS REACHED	CUB	320
C	IER=2	WORK SPACE FULL	CUB	330
C	IER=3	FURTHER SUBDIVISION OF TRIANGLES IMPOSSIBLE	CUB	340
C	IER=4	NO FURTHER IMPROVEMENT IN ACCURACY IS	CUB	350
C		POSSIBLE DUE TO ROUNDING ERRORS IN FUNCTION	CUB	360
C		VALUES	CUB	370
C	IER=5	NO FURTHER IMPROVEMENT IN ACCURACY IS	CUB	380
C		POSSIBLE BECAUSE SUBDIVISION DOES NOT	CUB	390
C		CHANGE THE ESTIMATED INTEGRAL. MACHINE	CUB	400
C		ACCURACY HAS PROBABLY BEEN REACHED BUT	CUB	410
C		THE ERROR ESTIMATE IS NOT SHARP ENOUGH.	CUB	420
C			CUB	430
C		CUBTRI IS DESIGNED TO BE CALLED REPEATEDLY WITHOUT WASTING	CUB	440
C		EARLIER WORK. THE PARAMETER NCALLS IS USED TO INDICATE TO	CUB	450
C		CUBTRI AT WHAT POINT TO RESTART, AND MUST BE RE-INITIALIZED	CUB	460
C		TO 0 WHEN A NEW INTEGRAL IS TO BE COMPUTED. AT LEAST ONE OF	CUB	470
C		THE PARAMETERS EPS, MCALLS AND NW MUST BE CHANGED BETWEEN	CUB	480
C		CALLS TO CUBTRI, ACCORDING TO THE RETURNED VALUE OF IER. NONE	CUB	490
C		OF THE OTHER PARAMETERS MAY BE CHANGED IF RESTARTING IS DONE.	CUB	500
C		IF IER=3 IS ENCOUNTERED, THERE PROBABLY IS A SINGULARITY	CUB	510
C		SOMEWHERE IN THE REGION. THE ERROR MESSAGE INDICATES THAT	CUB	520
C		FURTHER SUBDIVISION IS IMPOSSIBLE BECAUSE THE VERTICES OF THE	CUB	530
C		SMALLER TRIANGLES PRODUCED WILL BEGIN TO COALESCE TO THE	CUB	540
C		PRECISION OF THE COMPUTER. THIS SITUATION CAN USUALLY BE	CUB	550
C		RELIEVED BY SPECIFYING THE REGION IN SUCH A WAY THAT THE	CUB	560
C		SINGULARITY IS LOCATED AT THE THIRD VERTEX OF THE TRIANGLE.	CUB	570
C		IF IER=4 IS ENCOUNTERED, THE VALUE OF THE INTEGRAL CANNOT BE	CUB	580
C		IMPROVED ANY FURTHER. THE ONLY EXCEPTION TO THIS OCCURS WHEN A	CUB	590
C		FUNCTION WITH HIGHLY IRREGULAR BEHAVIOUR IS INTEGRATED (E.G.	CUB	600
C		FUNCTIONS WITH JUMP DISCONTINUITIES OR VERY HIGHLY OSCILLATORY	CUB	610
C		FUNCTIONS). IN SUCH A CASE THE USER CAN DISABLE THE ROUNDING	CUB	620
C		ERROR TEST BY REMOVING THE IF STATEMENT SHORTLY AFTER STATEMENT	CUB	630
C		NUMBER 70.	CUB	640
C			CUB	650
C				
C		SUBROUTINE CUBRUL(F, VEC, P, IDATA, RDATA)	CUB	10
C			CUB	20
C		BASIC CUBATURE RULE PAIR OVER A TRIANGLE	CUB	30
C			CUB	40
C		PARAMETERS	CUB	50
C		F - EXTERNAL FUNCTION - SEE COMMENTS TO CUBTRI	CUB	60
C		VEC- MATRIX OF BASE VECTORS AND ORIGIN (INPUT)	CUB	70
C		P - TRIANGLE DESCRIPTION VECTOR OF DIMENSION 6	CUB	80
C		P(1) - TRANSFORMED X COORDINATE OF ORIGIN VERTEX(INPUT)	CUB	90
C		P(2) - TRANSFORMED Y COORDINATE OF ORIGIN VERTEX(INPUT)	CUB	100
C		P(3) - DISTANCE OF OTHER VERTICES IN THE DIRECTIONS	CUB	110
C		OF THE BASE VECTORS (INPUT)	CUB	120
C		P(4) - LESS ACCURATE ESTIMATED INTEGRAL (OUTPUT)	CUB	130
C		P(5) - MORE ACCURATE ESTIMATED INTEGRAL (OUTPUT)	CUB	140
C		P(6) - ABS(P(5)-P(4)) (OUTPUT)	CUB	150
C			CUB	160
C		CUBRUL EVALUATES A LINEAR COMBINATION OF BASIC INTEGRATION	CUB	170
C		RULES HAVING D3 SYMMETRY. THE AREAL COORDINATES PERTAINING TO	CUB	180
C		THE J-TH RULE ARE STORED IN W(I,J), I=1,2,3. THE CORRESPONDING	CUB	190
C		WEIGHTS ARE W(4,J) AND W(5,J), WITH W(5,J) BELONGING TO THE	CUB	200
C		MORE ACCURATE FORMULA. IF W(1,J).EQ.W(2,J), THE INTEGRATION	CUB	210
C		POINT IS THE CENTROID, ELSE IF W(2,J).EQ.W(3,J), THE EVALUATION	CUB	220
C		POINTS ARE ON THE MEDIANS. IN BOTH CASES ADVANTAGE IS TAKEN OF	CUB	230
C		SYMMETRY TO AVOID REPEATING FUNCTION EVALUATIONS.	CUB	240
C			CUB	250
C		THE FOLLOWING DOUBLE PRECISION VARIABLES ARE USED TO AVOID	CUB	260
C		UNNECESSARY ROUNDING ERRORS IN FLOATING POINT ADDITION.	CUB	270
C		THEY MAY BE DECLARED SINGLE PRECISION IF DOUBLE PRECISION IS	CUB	280
C		NOT AVAILABLE AND FULL ACCURACY IS NOT NEEDED.	CUB	290
C		DOUBLE PRECISION A1, A2, S, SN, DZERO, DONE, DTHREE, DSIX	CUB	300
C		REAL AREA, ORIGIN(2), P(6), RDATA(1), TVEC(2,3), VEC(2,3), W(5,6)	CUB	310
C		INTEGER IDATA(1)	CUB	320
C			CUB	330
C		W CONTAINS POINTS AND WEIGHTS OF THE INTEGRATION FORMULAE	CUB	340
C		NQUAD - NUMBER OF BASIC RULES USED	CUB	350
C			CUB	360
C		THIS PARTICULAR RULE IS THE 19 POINT EXTENSION (DEGREE 8) OF	CUB	370

C	THE FAMILIAR 7 POINT RULE (DEGREE 5).	CUB 380
C		CUB 390
C	SIGMA=SQRT(7)	CUB 400
C	PHI=SQRT(15)	CUB 410
C	W(1,1),W(2,1),W(3,1) = 1/3	CUB 420
C	W(4,1) = 9/40	CUB 430
C	W(5,1) = 7137/62720 - 45*SIGMA/1568	CUB 440
C	W(1,2) = 3/7 + 2*PHI/21	CUB 450
C	W(2,2),W(3,2) = 2/7 - PHI/21	CUB 460
C	W(4,2) = 31/80 - PHI/400	CUB 470
C	W(5,2) = - 9301697/4695040 - 13517313*PHI/23475200	CUB 480
C	+ 764885*SIGMA/939008 + 198763*PHI*SIGMA/939008	CUB 490
C	W(*,3) = W(*,2) WITH PHI REPLACED BY -PHI	CUB 500
C	W(1,5) = 4/9 + PHI/9 + SIGMA/9 - SIGMA*PHI/45	CUB 510
C	W(2,5),W(3,5) = 5/18 - PHI/18 - SIGMA/18 + SIGMA*PHI/90	CUB 520
C	W(4,5) = 0	CUB 530
C	W(5,5) = 102791225/59157504 + 23876225*PHI/59157504	CUB 540
C	- 34500875*SIGMA/59157504 - 9914825*PHI*SIGMA/59157504	CUB 550
C	W(*,4) = W(*,5) WITH PHI REPLACED BY -PHI	CUB 560
C	W(1,6) = 4/9 + SIGMA/9	CUB 570
C	W(2,6) = W(2,4)	CUB 580
C	W(3,6) = W(2,5)	CUB 590
C	W(4,6) = 0	CUB 600
C	W(5,6) = 11075/8064 - 125*SIGMA/288	CUB 610
C		CUB 620

ALGORITHM 585

A Subroutine for the General Interpolation and Extrapolation Problems

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1. INTRODUCTION

Recently a generalization of Neville-Aitken algorithm for extrapolation has been found by Håvie [10]. This general extrapolation scheme has also been studied by Brezinski [3], who showed that it includes most of the convergence acceleration methods actually known. This algorithm has been called the E-algorithm, or the BH-protocol [15, Chaps. 10 and 11, pp. 175–209].

Of course, after that, the next step was to look for a general interpolation scheme for interpolation by a linear combination of functions forming a Chebyshev system. Such an algorithm was in fact found some years ago by Mühlbach [12, 13]. Except for the initializations, it is the same as the E-algorithm. This algorithm has been called the Mühlbach–Neville-Aitken algorithm (MNA-algorithm). The same algorithm can also be used for the general interpolation problem [4] as described, for example, by Davis [7].

The E-algorithm and its programming are described in Section 2. A FORTRAN subroutine is given with an application. Section 3 is devoted to the MNA-algorithm, while Section 4 deals with least squares approximation and extrapolation.

2. THE E-ALGORITHM

Let (S_n) be a sequence of real numbers and let us consider the sequence transformation defined by

$$E_k(S_n) = \frac{\begin{vmatrix} S_n & S_{n+1} & \cdots & S_{n+k} \\ g_1(n) & g_1(n+1) & \cdots & g_1(n+k) \\ g_k(n) & g_k(n+1) & \cdots & g_k(n+k) \end{vmatrix}}{\begin{vmatrix} 1 & 1 & \cdots & 1 \\ g_1(n) & g_1(n+1) & \cdots & g_1(n+k) \\ g_k(n) & g_k(n+1) & \cdots & g_k(n+k) \end{vmatrix}}$$

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where the g_i 's are given sequences such that the denominator in $E_k(S_n)$ does not vanish.

The main basic algebraic property of this transformation is that $\forall n \geq N, E_k(S_n) = S$ if the sequence (S_n) satisfies, $\forall n \geq N$

$$S_n = S + a_1 g_1(n) + \dots + a_k g_k(n).$$

This formalism is quite general since it includes most of the sequence transformations actually used to accelerate the convergence of the sequence (S_n) . For example, if $g_i(n) = \Delta S_{n+i-1}$, it is Shanks' transformation [14]; if $g_i(n) = x_n^{i-1} \Delta S_n$, it is Levin's method [11]; for $g_i(n) = x_n^i$ we recover Richardson extrapolation process; and for $g_i(n) = (\Delta S_n)^i$, it is Germain-Bonne's method [9]; and so forth.

From the numerical point of view, a recursive method to avoid the computation of the determinants involved in $E_k(S_n)$ is needed. Such algorithms exist for the particular cases quoted above: Wynn's ϵ -algorithm [16] for Shanks' transformation, Neville-Aitken's scheme for Richardson process and Germain-Bonne's transformation, some special devices [5] for Levin's method, and so on. The corresponding FORTRAN subroutines can be found in [1]. Thus we have at least as many different algorithms as sequence transformations.

Recently a recursive algorithm working in the general case has been obtained by Håvie [10] and Brezinski [3]. This algorithm, called the E-algorithm, is as follows:

$$\begin{aligned} \text{Initializations} \quad E_0^{(n)} &= S_n, & n &= 0, 1, \dots \\ g_{0,i}^{(n)} &= g_i(n), & i &= 1, 2, \dots \quad \text{and} \quad n = 0, 1, \dots \end{aligned}$$

Then, for $k = 1, 2, \dots$ and $n = 0, 1, \dots$

$$E_k^{(n)} = E_{k-1}^{(n)} + g_{k-1,k}^{(n)} \frac{E_{k-1}^{(n)} - E_{k-1}^{(n+1)}}{g_{k-1,k}^{(n+1)} - g_{k-1,k}^{(n)}} \quad (2.1)$$

$$g_{k,i}^{(n)} = g_{k-1,i}^{(n)} + g_{k-1,k}^{(n)} \frac{g_{k-1,i}^{(n)} - g_{k-1,i}^{(n+1)}}{g_{k-1,k}^{(n+1)} - g_{k-1,k}^{(n)}}, \quad i = k+1, k+2, \dots \quad (2.2)$$

It can very easily be proved, by using Sylvester's determinantal identity [8, Vol. 1, p. 32] that

$$E_k^{(n)} = E_k(S_n).$$

Usually the quantities $E_k^{(n)}$ are displayed in a two-dimensional array, called the E-array, as follows:

$$\begin{array}{cccc} E_0^{(0)} & = & S_0 & \\ E_0^{(1)} & = & S_1 & E_1^{(0)} \\ E_0^{(2)} & = & S_2 & E_1^{(1)} & E_2^{(0)} \\ E_0^{(3)} & = & S_3 & E_1^{(2)} & E_2^{(1)} & E_3^{(0)} \\ & & & \vdots & \vdots & \vdots \\ & & & \vdots & \vdots & \vdots \end{array}$$

The $g_{k,i}^{(n)}$'s can also be placed in similar arrays, the g_i -arrays, for a fixed value of the index i , where k indicates the column starting from zero. It is easy to check that $g_{k,i}^{(n)} = 0$ for $k \geq i$, and thus the g_i -array only contains the columns 0 to $i-1$.

There are two methods for using, and thus for programming, convergence acceleration methods. The first one is to use them a posteriori, that is, to compute a fixed number of terms of the sequence (S_n) and then to apply the algorithm by computing each column of the array from the preceding one. This use is very easy to program, but it has the main disadvantage that all the computations must be started again from the beginning if one wants to add new terms to the initial sequence (S_n) .

The second use, which is much more convenient in practice, is to use the algorithm "in parallel" with the computation of the successive terms of the original sequence (S_n) : After the computation of each new term of the sequence

Let us give an example of the use of this subroutine. It corresponds to the choice $g_i(n) = \Delta S_{n+i-1}$ (Shanks' transformation) for the sequence $S_n = 1/(n+1)$. For this sequence it is known that $E_k^{(0)} = \epsilon_{2k}^{(0)} = 1/(k+1)^2$.

```

C      PROGRAM TESTD (OUTPUT,TAPE6=OUTPUT)                MAN  10
C      TEST DRIVER FOR THE EXTRAP SUBROUTINE              MAN  20
C      IN THE E-ALGORITHM                                MAN  30
C                                                        MAN  40
C      INTEGER I, MAXCOL, INIT, INFO, K, M, MPI, NI      MAN  50
C      DOUBLE PRECISION KPI, SK, RESULT, S, GINIT, EPS, T MAN  60
C      DIMENSION GINIT(50)                               MAN  70
C                                                        MAN  80
C      DATA NOUT /6/, INIT /0/, MAXCOL /5/, EPS /1.D-30/ MAN  90
C                                                        MAN 100
C      DO 40 K=1,20                                       MAN 110
C                                                        MAN 120
C      COMPUTATION OF S                                   MAN 130
C      K-1                                               MAN 140
C                                                        MAN 150
C      T = K                                             MAN 160
C      SK = 1.D0/T                                       MAN 170
C                                                        MAN 180
C      COMPUTATION OF G (K-1)                            MAN 190
C      I                                                 MAN 200
C                                                        MAN 210
C      M = MAX0(1,K-2)                                   MAN 220
C      DO 10 I=1,M                                       MAN 230
C      KPI = K + I                                       MAN 240
C      GINIT(I) = -1.D0/(KPI-1.D0)/KPI                  MAN 250
10  CONTINUE                                           MAN 260
C      IF (K.LT.3 .OR. K.GT.MAXCOL+1) GO TO 30          MAN 270
C                                                        MAN 280
C      COMPUTATION OF G (I-1)                            MAN 290
C      K-1                                               MAN 300
C                                                        MAN 310
C      DO 20 I=1,K                                       MAN 320
C      KPI = K + I                                       MAN 330
C      MPI = M + I                                       MAN 340
C      GINIT(MPI) = -1.D0/(KPI-2.D0)/(KPI-1.D0)        MAN 350
20  CONTINUE                                           MAN 360
30  CALL EXTRAP(INIT, EPS, MAXCOL, SK, GINIT, RESULT, INFO) MAN 370
C      IF (INFO.LT.0) WRITE (NOUT,99998)                MAN 380
C      IF (INFO.EQ.1) WRITE (NOUT,99997)                MAN 390
C      T = K                                             MAN 400
C      S = 1.D0/T**2                                     MAN 410
C      T = (MAXCOL+1)*K                                  MAN 420
C      IF (K.GT.MAXCOL+1) S = 1.D0/T                    MAN 430
C      WRITE (NOUT,99999) RESULT, S                     MAN 440
40  CONTINUE                                           MAN 450
C      STOP                                             MAN 460
99999 FORMAT (3D26.12)                                  MAN 470
99998 FORMAT (46H DIMENSIONS TOO SMALL IN THE SUBROUTINE EXTRAP) MAN 480
99997 FORMAT (42H DIVISION BY ZERO IN THE SUBROUTINE EXTRAP)  MAN 490
C      END                                             MAN 500

```

This program gives the results shown in Table I on a computer with precision 10^{-18} (DEC 10). The numbers printed are $E_0^{(0)}, \dots, E_5^{(0)}, E_5^{(1)}, \dots, E_5^{(14)}$, where the exact values are given by $E_k^{(n)} = 1/((n+k+1)(k+1))$. The E-algorithm can be quite sensitive to the propagation of rounding errors. Several possibilities exist for their control. The first one is to compute the "amplification factors" defined by Håvie [9] which can be used to obtain an upper bound for the amplification of the initial errors and which can be considered as a measurement of the numerical stability of the algorithm. The second possibility consists in using the particular rules given in [4] although they are quite difficult to program and will lead to a much longer program. The third possibility is to use an exact arithmetic such as p-adic or rational arithmetic. Such an arithmetic is actually under consideration.

A quite important practical problem when using convergence acceleration methods is that of the termination criteria. For some special sequences and some particular methods such a criterion can be obtained (see, for example, the discussion on totally monotonic and totally oscillating sequences and their treatment by Shanks's transformation in [1]). In the general case some theorems exist [3, Theorem 3], but in practice the computations are stopped when two successive elements of the same column or the same diagonal of the E-array are nearly equal.

To end this section let us mention that the relationships of the E-algorithm, as they are written in [3], are numerically unstable. These relations must be written

Table I

Computed values	Exact values
0.10000000000D+01	0.10000000000D+01
0.25000000000D+00	0.25000000000D+00
0.11111111111D+00	0.11111111111D+00
0.62500000000D-01	0.62500000000D-01
0.40000000000D-01	0.40000000000D-01
0.27777777778D-01	0.27777777778D-01
0.238095238096D-01	0.238095238095D-01
0.20833333332D-01	0.20833333333D-01
0.185185185190D-01	0.185185185185D-01
0.166666666650D-01	0.16666666667D-01
0.151515151566D-01	0.151515151515D-01
0.138888888769D-01	0.138888888889D-01
0.128205128394D-01	0.128205128205D-01
0.119047618869D-01	0.119047619048D-01
0.111111111178D-01	0.11111111111D-01
0.104166666706D-01	0.10416666667D-01
0.980392156462D-02	0.980392156863D-02
0.925925926345D-02	0.925925925926D-02
0.877192979360D-02	0.877192982456D-02
0.833333342965D-02	0.83333333333D-02

as they are in (2.1) and (2.2) and in the subroutine. In that way they will be much more stable. For example, if the unstable version of these rules is used, then, for the preceding example, we shall obtain for the last entry a result that is less accurate by three digits.

3. THE MÜHLBACH-NEVILLE-AITKEN ALGORITHM

Let $\{f_i\}_{i \geq 0}$ be a family of linearly independent elements of a real vector space E . $\{f_i\}$ is assumed to be a complete Chebyshev system with respect to the linear functionals $\{L_j\}$, which means that all the Gram determinants $|L_j(f_i)|_{i,j=0,\dots,k}$ are not zero whenever L_0, \dots, L_k are linearly independent.

The problem is to construct a generalized polynomial

$$P_k = a_0 f_0 + \dots + a_k f_k$$

such that

$$L_i(P_k) = w_i, \quad i = 0, \dots, k,$$

where the w_i 's are given numbers not all zero.

More generally one wants to construct

$$P_k^{(n)} = a_0 f_0 + \dots + a_k f_k$$

such that

$$L_i(P_k^{(n)}) = w_i, \quad i = n, \dots, n+k.$$

This is the general interpolation problem as described, for example, by Davis [7].

Usually on a computer (and also often in numerical analysis) one deals with numbers and not with elements of a vector space. This means that, in practice, one has to compute the numerical values $L(P_k^{(n)})$ where L is a linear functional on E . (If E is a vector space of functions, then, for example, $L(f)$ is the value of the function f at some prescribed point x .)

Mühlbach [12, 13] showed that these values can be recursively computed by an algorithm that generalizes the Neville-Aitken scheme. For this reason it has been called the Mühlbach-Neville-Aitken algorithm (briefly, the MNA-algorithm). A very simple proof of the MNA-algorithm has been given in [4] using Sylvester's determinantal identity. Except for the initializations, it is similar to the E-algorithm.

C	(K-1)	MAN	570
C	COMPUTATION OF F	MAN	580
C	ϕ, I	MAN	590
C		MAN	600
	DO 10 I=1,M	MAN	610
	GINIT(I) = A(K,I+1)*F(1)/A(K,I) - F(I+1)	MAN	620
10	CONTINUE	MAN	630
	IF (K.LT.3) GO TO 30	MAN	640
C		MAN	650
C	(I-1)	MAN	660
C	COMPUTATION OF F	MAN	670
C	$\phi, K-1$	MAN	680
C		MAN	690
	DO 20 I=1,K	MAN	700
	MPI = M + I	MAN	710
	GINIT(MPI) = A(I,K)*F(1)/A(I,I) - F(K)	MAN	720
20	CONTINUE	MAN	730
30	CALL EXTRAP(INIT, EPS, N-1, SK, GINIT, RESULT, INFO)	MAN	740
40	CONTINUE	MAN	750
	IF (INFO.LT.0) WRITE (NOUT,99998)	MAN	760
	IF (INFO.EQ.1) WRITE (NOUT,99997)	MAN	770
	P = 3.D0*F(1) + 4.D0*F(2) - F(3) + F(4)	MAN	780
	WRITE (NOUT,99999) X, P, RESULT	MAN	790
50	CONTINUE	MAN	800
	STOP	MAN	810
99999	FORMAT (3D26.12)	MAN	820
99998	FORMAT (46H DIMENSIONS TOO SMALL IN THE SUBROUTINE EXTRAP)	MAN	830
99997	FORMAT (42H DIVISION BY ZERO IN THE SUBROUTINE EXTRAP)	MAN	840
	END	MAN	850

Table II

x	$L(f)$	$L(P_4^{(0)})$
0.20000000000000D+00	0.16014316699801D+01	0.16014316699801D+01
0.40000000000000D+00	0.24923427754674D+01	0.24923427754674D+01
0.60000000000000D+00	0.30341791780962D+01	0.30341791780962D+01
0.80000000000000D+00	0.33542677046355D+01	0.33542677046355D+01
0.10000000000000D+01	0.34695222627459D+01	0.34695222627459D+01
0.12000000000000D+01	0.33740605363995D+01	0.33740605363995D+01
0.14000000000000D+01	0.30580620483999D+01	0.30580620483999D+01
0.16000000000000D+01	0.25139875953957D+01	0.25139875953957D+01
0.18000000000000D+01	0.17390738530797D+01	0.17390738530797D+01
0.20000000000000D+01	0.73634273757251D+00	0.73634273757251D+00

Of course, $L(P_4^{(0)}) = L(f)$ for all x , as can be checked from the numerical results shown in Table II.

The MNA-algorithm can also be used to implement another convergence acceleration process due to Germain-Bonne [9, pp. 33-43] that is based on generalized inverse interpolation.

4. LEAST SQUARES APPROXIMATION AND EXTRAPOLATION

Let E be an inner product space. The least squares approximation of $f \in E$ by a combination of the given independent elements f_0, \dots, f_k of E consists in determining $\alpha_0, \dots, \alpha_k$ such that

$$\|f - (\alpha_0 f_0 + \dots + \alpha_k f_k)\|$$

be minimum where $\|\cdot\|^2 = (\cdot, \cdot)$. It is well known (e.g., see [7, p. 182]) that the solution $F_k = \alpha_0 f_0 + \dots + \alpha_k f_k$ of this minimization problem is given by a ratio of determinants completely similar to the one defining $P_k^{(0)}$ in the general interpolation problem; one only has to replace $L_i(f_j)$ by (f_i, f_j) and w_i by (f_i, f) . If L is a linear functional on E , then the sequence $(L(F_k))$ can be obtained via the MNA algorithm with the initializations

$$P_0^{(n)} = (f_n, f) \frac{L(f_0)}{(f_n, f_0)}$$

$$f_{0,i}^{(n)} = (f_n, f_i) \frac{L(f_0)}{(f_n, f_0)} - L(f_i).$$

We shall obtain

$$P_k^{(0)} = L(F_k).$$

A special case of this general problem is to find F_k , which satisfies the system of linear equations

$$L_i(F_k) = w_i, \quad i = 0, \dots, m \geq k$$

in the least squares sense. The solution of this problem can be obtained by the MNA-algorithm where the scalar product is defined by

$$(f, g) = \sum_{i=0}^m L_i(f)L_i(g).$$

Another special case is the biconjugate gradient method for solving systems of linear equations whose successive iterates are still given by a similar ratio of determinants [2, p. 87] and which can therefore be computed by the MNA-algorithm.

The subroutine EXTRAP can be used for implementing all these least squares approximation problems. This subroutine can also be used for a convergence acceleration process due to Cordellier [6] that is based on least squares. Let us explain this process.

The E-transformation of sequences described in Section 2 was obtained by assuming that the given sequence (S_n) satisfied

$$S_{n+i} = S + a_1 g_1(n+i) + \dots + a_k g_k(n+i)$$

for $i = 0, \dots, k$ and then solving this system for the unknown S . Cordellier's idea consists of writing this system for $i = 0, \dots, m \geq k$ and then solving it in the least squares sense for the unknown S . This least squares solution can be obtained by the E-algorithm as we now see. We set

$$\begin{aligned} g_i &= (g_i(n); \dots; g_i(n+m))^T, \quad i \geq 1 \\ g_0 &= (1; \dots; 1)^T \\ V &= (S_n; \dots; S_{n+m})^T. \end{aligned}$$

Thus the system becomes

$$\begin{pmatrix} 1 & \frac{(g_0, g_1)}{(g_0, g_0)} & \dots & \frac{(g_0, g_k)}{(g_0, g_0)} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & \frac{(g_k, g_1)}{(g_k, g_0)} & \dots & \frac{(g_k, g_k)}{(g_k, g_0)} \end{pmatrix} \begin{pmatrix} S \\ a_1 \\ \vdots \\ a_k \end{pmatrix} = \begin{pmatrix} \frac{(g_0, V)}{(g_0, g_0)} \\ \vdots \\ \frac{(g_k, V)}{(g_k, g_0)} \end{pmatrix}$$

The value of S thus obtained will depend on m , k , and n , and we denote it by ${}_m E_k^{(n)}$. Since the preceding system has the same form as the system implicitly solved by the E-algorithm, we can use it to obtain the sequence ${}_m E_k^{(n)}$ for $k = 0, \dots, m$ and n and m fixed. In this case the initializations must be

$$\begin{aligned} E_0^{(n)} &= \frac{(g_n, V)}{(g_n, g_0)} \\ g_{0,i}^{(n)} &= \frac{(g_n, g_i)}{(g_n, g_0)}, \quad i \geq 1 \end{aligned}$$

and we shall get

$$E_k^{(0)} = {}_m E_k^{(n)}.$$

If the value of n or m is modified, then the definition of the scalar product is also modified and all the computations have to be started again from the beginning. For $k = m$ the least squares extrapolated value will be identical to the value obtained by the E-transformation as described in Section 2. This is also

Table III

k	${}_4E_k^{(0)}$ (Perturbed sequence and perturbed g_i)	R.E.	${}_4E_k^{(0)}$ (Perturbed sequence and unperturbed g_i)	R.E.	${}_4E_k^{(0)}$ (Unperturbed sequence)
	0	0.4567459	1.7	0.4567459	1.7
1	0.1782071	2.3	0.1782336	3.8	0.1781660
2	0.0937578	6.4	0.0938553	4.0	0.0938178
3	0.0581327	32.4	0.0583042	3.0	0.0583216
4	0.0398319	42.0	0.0398976	25.6	0.0399999

true for the value obtained by least squares approximation that will be equal to the value obtained by interpolation when $m = k$.

The basic advantage of least squares extrapolation is to be less sensitive to small perturbations in the sequence (S_n) to be accelerated.

Let us take the same example as in Section 2 and let us add a perturbation of magnitude less than 10^{-4} to the sequence (S_n) .

Since the sequences (g_i) depend on (S_n) , a perturbation of (S_n) will also modify the (g_i) . Thus two different experiments have been conducted: with the (g_i) automatically modified by the perturbation of (S_n) and with unmodified (g_i) . Columns R.E. in Table III contain the values of the relative errors with respect to the unperturbed case multiplied by 10000. For $n = 0$ and $m = 4$ we get the results also shown in the table. As we can verify, the ${}_4E_k^{(0)}$'s ($k < 4$) are less sensitive to the perturbation than ${}_4E_4^{(0)} = E_4^{(0)} = 0.04$. It must be noticed that the results obtained are very sensitive to the precision on the computation of the scalar products appearing in the initializations.

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ALGORITHM 586

ITPACK 2C: A FORTRAN Package for Solving Large Sparse Linear Systems by Adaptive Accelerated Iterative Methods

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1. INTRODUCTION

ITPACK 2C is a collection of seven FORTRAN subroutines for solving large sparse linear systems by adaptive accelerated iterative algorithms. Basic iterative procedures, such as the Jacobi method, the successive overrelaxation method, the symmetric successive overrelaxation method, and the RS method for the reduced system are combined, where possible, with acceleration procedures, such as Chebyshev (semi-iteration) and conjugate gradient, for rapid convergence. Automatic selection of the acceleration parameters and the use of accurate stopping criteria are major features of this software package. While the ITPACK routines can be called with any linear system containing positive diagonal elements, they are the most successful in solving systems with symmetric positive definite or mildly nonsymmetric coefficient matrices.

For several years, we have been involved with the development and use of research-oriented programs using iterative algorithms for solving large sparse

$$Au = b$$

with positive diagonal elements. One solves for the N component unknown vector u given the $N \times N$ nonsingular coefficient matrix A and the N component right-hand side vector b . The current ITPACK software package of subroutines, version 2C, provides for the use of seven alternative iterative procedures. While these subroutines are not designed as production software, they should successfully handle industrial problems of moderate size, that is, ones that fit in high-speed

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memory. This package is written in 1966 American National Standard FORTRAN code. It has been tested over a wide variety of computer systems using various FORTRAN compilers, including one which is FORTRAN-77 compatible (see Acknowledgments).

The seven iterative solution modules are based on several basic iterative procedures, such as the Jacobi method, the successive overrelaxation (SOR) method, the symmetric SOR (SSOR) method, and the RS method for the reduced system. With the exception of SOR, the convergence of these basic methods are accelerated by Chebyshev (semi-iteration, SI) or conjugate gradient (CG) acceleration. All methods are available with adaptive parameter estimation and automatic stopping tests. When using the RS method it is required that the linear system be reordered into a "red-black" system [6, 13].¹ A switch to compute, if possible, the red-black indexing, permute the linear system, and permute associated vectors is provided.

The successful convergence of iterative methods may be dependent on conditions that are difficult to determine in advance. For example, determining whether the coefficient matrix is positive definite can be as costly to check as solving the system. On the other hand, some conditions affecting convergence, such as positive diagonal elements, diagonal dominance, and symmetry are relatively easy to verify. For some applications, the theory to guarantee the convergence of an iterative method may not exist. The algorithms in ITPACK have been tested most extensively for linear systems arising from elliptic partial differential equations. The routines can be applied formally to any linear system which fits in high-speed memory. However, rapid convergence, and indeed convergence itself cannot be guaranteed unless the matrix of the system is symmetric and positive definite. Success can be expected, though not guaranteed, for mildly nonsymmetric systems. In other words, iterative methods may not converge when applied to systems with coefficient matrices which are completely general with no special properties.

This article discusses the usage of ITPACK and gives a few test results. The description of the iterative methods is given in [4]. The underlying theory on which the iterative algorithms are based is described in [6]. A survey of the iterative methods in ITPACK is presented in [12].

Throughout this paper, we adopt notation such as **SOR()** when referring to a subroutine and **A(*)** for a single-dimensioned array. The residual vector is $b - Au$ for the linear system $Au = b$ and the pseudo-residual vector is $Gu^{(n)} + k - u^{(n)}$ for a basic iterative method of the form $u^{(n+1)} = Gu^{(n)} + k$. The smallest and largest eigenvalues of the iteration matrix G are denoted $m(G)$ and $M(G)$, respectively.

2. SPARSE MATRIX STORAGE

The sparse storage scheme used in ITPACK is a common one. It is a row-wise representation of the nonzero entries in the coefficient matrix of the linear system. For a nonsymmetric coefficient matrix, all of the nonzero values in each row are stored in a contiguous block of data in a real-valued array **A(*)**. If the matrix is symmetric, computer memory can be saved by storing only the nonzero entries in each row on and above the main diagonal. For either nonsymmetric or symmetric sparse storage, associated column numbers are stored in an integer-valued array **JA(*)** such that **JA(K)** is the column number for the value **A(K)**.

¹ In this ordering, the components of the unknown vector u are considered as either "red" or "black". A "red-black ordering" is any ordering such that every black unknown follows all of the red unknowns. This ordering of unknowns leads to a 2×2 "red-black partitioning" of the coefficient matrix; that is, a matrix of the form

$$\begin{bmatrix} A_{1,1} & A_{1,2} \\ A_{2,1} & A_{2,2} \end{bmatrix}$$

with diagonal submatrices $A_{1,1}$ and $A_{2,2}$. The original linear system may require rearrangement in order to arrive at this form.

A mapping vector \mathbf{IA}^* is used to denote the starting locations of each of the contiguous blocks. The beginning of the linear block for row \mathbf{I} is given by $\mathbf{IA}(\mathbf{I})$, the end by $\mathbf{IA}(\mathbf{I} + 1) - 1$, and its length by $\mathbf{IA}(\mathbf{I} + 1) - \mathbf{IA}(\mathbf{I})$. Thus, \mathbf{IA}^* will contain $\mathbf{N} + 1$ elements to accommodate a linear system of order \mathbf{N} . The entries for each row may be stored in any order in the contiguous block for that row.

For example, the coefficient matrix

$$\begin{bmatrix} 11. & 0. & 0. & 14. & 15. \\ 0. & 22. & 0. & 0. & 0. \\ 0. & 0. & 33. & 0. & 0. \\ 14. & 0. & 0. & 44. & 45. \\ 15. & 0. & 0. & 45. & 55. \end{bmatrix}$$

would be represented in nonsymmetric sparse storage as

$$\mathbf{A}^* = [11., 14., 15., 22., 33., 14., 44., 45., 15., 45., 55.]$$

$$\mathbf{JA}^* = [1, 4, 5, 2, 3, 1, 4, 5, 1, 4, 5]$$

$$\mathbf{IA}^* = [1, 4, 5, 6, 9, 12]$$

and in symmetric sparse storage as

$$\mathbf{A}^* = [11., 14., 15., 22., 33., 44., 45., 55.]$$

$$\mathbf{JA}^* = [1, 4, 5, 2, 3, 4, 5, 5]$$

$$\mathbf{IA}^* = [1, 4, 5, 6, 8, 9]$$

3. USAGE

The user is expected to provide the coefficient matrix and the right-hand side of the linear system to be solved. The data structure for the matrix of the system is either the symmetric or nonsymmetric sparse storage format described in the previous section. An initial guess for the solution should be provided, if one is known; otherwise, it can be set to all zero values. A series of approximations for the solution are generated iteratively until the convergence criterion is satisfied. The algorithms are performed in two work-space arrays and some control over the algorithmic procedure can be obtained from switches in two parameter arrays.

There are seven main subroutines in ITPACK, each corresponding to an iterative method. They are

<u>SUBROUTINE</u>	<u>METHOD</u>
JCG()	Jacobi Conjugate Gradient
JSI()	Jacobi Semi-iteration
SOR()	Successive Overrelaxation
SSORCG()	Symmetric SOR Conjugate Gradient
SSORSI()	Symmetric SOR Semi-iteration
RSCG()	Reduced System Conjugate Gradient
RSSI()	Reduced System Semi-iteration

and the calling sequence is

CALL (method) (N, IA, JA, A, RHS, U, IWKSP, NW, WKSP, IPARM, RPARM, IER)

where the parameters are defined in the following. Here "input" means that the subroutine expects the user to provide the necessary input data, and "output" means that the routine passes back information in the variable or array indicated. All parameters are linear arrays, except variables \mathbf{N} , \mathbf{NW} , and \mathbf{IER} . Moreover, all parameters may be altered by the subroutine call, except variables \mathbf{N} and \mathbf{NW} . (See Section 7 for additional details.)

PARAMETER	DESCRIPTION
N	Order of the linear system. [integer; input]
IA(*)	Vector of length $N + 1$ used in the sparse-matrix storage format. Contains the row pointers into JA(*) and A(*) . [integer array; input]
JA(*)	Vector of length NZ (defined in A(*) below) used in the sparse-matrix storage format. Contains the column numbers for the corresponding entries in A(*) . [integer array; input]
A(*)	Vector of length NZ used in the sparse-matrix storage format. Contains the nonzero entries of the coefficient matrix with positive diagonal elements. (NZ is the number of nonzero entries in the upper triangular part of the coefficient matrix when symmetric storage is used and is the total number of nonzeros when nonsymmetric storage is used.) [real array; input]
RHS(*)	Vector of length N containing the right-hand side of the linear system. [real array; input]
U(*)	Vector of length N containing the initial guess to the solution of the linear system on input and the latest approximate solution on output. [real array; input/output]
IWKSP(*)	Vector of length $3 * N$ used for integer work space. When reindexing for red-black ordering, the first N locations contain on output the permutation vector for the red-black indexing, the next N locations contain its inverse, and the last N are used for integer work space. ² [integer array; output]
NW	On input, NW is the available length for WKSP(*) . On output, IPARM(8) is the actual amount used (or needed). [integer; input]
WKSP(*)	Vector used for real working space whose length depends on the iterative method being used. Must be at least NW entries long. (See table near end of this section for required amount of work space for each method.) [real array]
IPARM(*)	Vector of length 12 used to initialize various integer and logical parameters. Default values may be set by calling subroutine DFAULT() described below. On output, IPARM(*) contains the values of the parameters that were changed. (Further details to follow.) [integer array; input/output]
RPARAM(*)	Vector of length 12 used to initialize various real parameters on input. Default values may be set by calling subroutine DFAULT() described below. On output, RPARAM(*) contains the final values of the parameters that were changed. (Further details given later in this section.) [real array; input/output]
IER	Error flag which is set to zero for normal convergence and to a nonzero integer when an error condition is present. (See table at end of section for meaning of nonzero values.) [integer; output]

The user may supply nondefault values for selected quantities in **IPARM(*)** and **RPARAM(*)** by first executing

CALL DFAULT (IPARM, RPARAM)

and then assigning the appropriate nondefault values before calling a solution module of **ITPACK**.

² For the red-black ordering, the *I*th entry of a permutation array **P()** indicates the position **J** into which the *I*th unknown of the original system is being mapped; that is, if **P(I) = J**, then unknown **I** is mapped into position **J**. The *J*th entry of an inverse permutation array **IP()** indicates the position **I** into which the *J*th unknown of the permuted system must be mapped to regain the original ordering, that is, **IP(J) = I**.

The iterative algorithms used in ITPACK are quite complicated and some knowledge of iterative methods is necessary to completely understand them. The interested reader should consult [4] and [6] for details. Important variables in this package which may change adaptively are **CME** (estimate of largest eigenvalue of the Jacobi matrix), **SME** (estimate of the smallest eigenvalue of the Jacobi matrix), **OMEGA** (overrelaxation parameter for the SOR and SSOR methods), **SPECR** (estimate of the spectral radius of the SSOR matrix), **BETAB** (estimate for the spectral radius of the matrix LU , where L and U are strictly lower and upper triangular matrices, respectively, such that the Jacobi matrix $B = L + U$).

The integer array **IPARM(*)** and real array **RPARAM(*)** allow the user to control certain parameters which affect the performance of the iterative algorithms. Furthermore, these arrays allow the updated parameters from the automatic adaptive procedures to be communicated back to the user. The entries in **IPARM(*)** and **RPARAM(*)** are

<u>POSITION</u>	<u>DEFAULT</u>	<u>MEANING OR USAGE</u>
IPARM(1)	100	ITMAX: Maximum number of iterations allowed. Reset on output to the number of iterations performed.
IPARM(2)	0	LEVEL: Control level of output. Each higher value provides additional information. [<0: no output on unit IPARM(4) ; 0: fatal error messages only; 1: warning messages and minimum output; 2: reasonable summary (progress of algorithm); 3: parameter values and informative comments; 4: approximate solution after each iteration (primarily useful for debugging); 5: original system]
IPARM(3)	0	IRESET: Communication switch. [0: implies certain values of IPARM(*) and RPARAM(*) will be overwritten to communicate parameters back to the user; otherwise: only IPARM(1) and IPARM(8) will be reset.]
IPARM(4)	6	NOOUT: output unit number.
IPARM(5)	0	ISYM: Sparse-storage format switch. [0: symmetric sparse storage; 1: nonsymmetric sparse storage]
IPARM(6)	1	IADAPT: Adaptive switch. Determines whether certain parameters have been set by user or should be computed automatically in either a fully or partially adaptive sense. [0: fixed iterative parameters used for SME , CME , OMEGA , SPECR , and BETAB (nonadaptive); 1: fully adaptive procedures used for all parameters; 2: (SSOR methods only) SPECR determined adaptively and CME , BETAB , and OMEGA fixed; 3: (SSOR methods only) BETAB fixed and all other parameters determined adaptively]

(See [4, 6] for details and **RPARAM(I)**, $I = 2, 3, 5, 6, 7$, for **CME**, **SME**, **OMEGA**, **SPECR**, **BETAB**, respectively. These parameters are

IPARM(7)	1	<p>set by subroutine DFAULT() or by the user.)</p> <p>ICASE: Adaptive procedure case switch for JSI and SSOR methods. There are two strategies, called Case I and Case II, for doing the adaptive procedure. The choice of which case to select corresponds to knowledge of the eigenvalues of the Jacobi matrix B and their estimates.</p> <p>[($\neq 2$) Case I: fixed $\mathbf{SME} \leq m(B)$ (general case); $(= 2)$ Case II: use when it is known that $m(B) \leq M(B)$]</p> <p>The case switch determines how the estimates for SME and CME are recomputed adaptively. In Case I, SME is fixed throughout and should be less than or equal to $m(B)$. In Case II, SME is set to $-\mathbf{CME}$, which may adaptively change. As far as the adaptive procedure is concerned, Case I is the most general case and should be specified in the absence of specific knowledge of the relationship between the eigenvalues and their estimates. An example when Case II is appropriate occurs when the Jacobi matrix has Property A, since $m(B) = -M(B)$.³ Also, if A is an L-matrix, then for the Jacobi matrix, we have $m(B) \leq M(B)$ and SME is always $-\mathbf{CME}$ (Case II).⁴ Selecting the correct case may increase the rate of convergence of the iterative method. (See [6] for additional discussion on Cases I and II. Also, see RPARAM(I), $I = 2, 3$ for CME, SME, respectively.)</p>
IPARM(8)	0	<p>NWKSP: Amount of work space used. Used for output only. If ITMAX is set to a value just over the actual number of iterations necessary for convergence, the amount of memory for WKSP(*) can be reduced to just over the value returned here. This may be done when rerunning a problem, for example.</p>
IPARM(9)	-1	<p>NB: Red-black ordering switch. On output, if reindexing is done, NB is set to the order of the black subsystem.</p> <p>[For RS methods, <0: compute red-black indexing and permute system; ≥ 0: skip indexing—system already in red-black form; For other methods, <0: skip indexing—system already in desired form; ≥ 0: compute red-black indexing and permute system]</p>

³ A matrix has Property A if and only if it is a diagonal matrix, or else there exists a rearrangement of the rows and corresponding columns of the matrix which corresponds to a red-black partitioning.

⁴ An L -matrix has positive diagonal elements and nonpositive off-diagonal elements.

		A negative integer value for IPARM(9) causes the equations to be handled in the most general way appropriate for the solution method being used. For methods other than RS methods this is the "natural order," while for RS methods it is the "red-black order." A nonnegative value produces a red-black permutation for all methods except for the RS methods, which are assumed to be in red-black order with the order of the black subsystem NB given. If reindexing is performed, IPARM(9) will contain the order of the black subsystem on output.
IPARM(10)	0	IREMOVE : Switch for effectively removing rows and columns when the diagonal entry is extremely large compared to the nonzero off-diagonal entries in that row. (See RPARM(8) for additional details.) [0: not done; otherwise: test done]
IPARM(11)	0	ITIME : Timing switch. [0: time method; otherwise: not done]
IPARM(12)	0	IDGTS : Error analysis switch. An analysis of final computed solution to determine accuracy. [<0: skip error analysis 0: compute DIGIT1 and DIGIT2 and store in RPARM(I) , I = 11, 12 , respectively; 1: print DIGIT1 and DIGIT2 ; 2: print final approximate solution vector; 3: print final approximate residual vector; 4: print both solution and residual vectors; otherwise: no printing] (If LEVEL \leq 0, no printing is done. See RPARM(I) , I = 11, 12 , for details on DIGIT1 and DIGIT2 .)

<u>ENTRY</u> <u>DEFAULT</u>	<u>DESCRIPTION</u>
RPARM(1) 5.E-6	ZETA : Stopping criterion or approximate relative accuracy desired in the final computer solution. If the method does not converge in IPARM(1) iterations, RPARM(1) is reset to an estimate of the relative accuracy achieved. The stopping criterion is a test of whether ZETA is greater than the ratio of the Euclidean norm of the pseudo-residual vector and the norm of the current iteration vector times a constant involving an eigenvalue estimate. (See [4, 6] for details.)
RPARM(2) 0.	CME : Estimate of largest eigenvalue of Jacobi matrix. Changes to new estimate if adaptive procedure is used. $CME \leq M(B)$.
RPARM(3) 0.	SME : Estimate of smallest eigenvalue of Jacobi matrix for JSI method. In Case I, SME is fixed throughout at a value $\leq m(B)$. In Case II, SME is always set to $-CME$ with CME changing in the adaptive procedure. (See IPARM(7) for definitions of Cases I and II.)

- RPARM(4) .75** **FF:** Adaptive procedure damping factor. Values in the interval $(0., 1.]$ with 1. causing the most frequent parameter changes when the fully adaptive switch **IPARM(6) = 1** is used.
- RPARM(5) 1.** **OMEGA:** Overrelaxation parameter for SOR and SSOR methods. If method is fully adaptive, **OMEGA** changes.
- RPARM(6) 0.** **SPECR:** Estimated spectral radius for SSOR matrix. If method is adaptive, **SPECR** changes.
- RPARM(7) .25** **BETAB:** Estimate for the spectral radius of matrix LU used in SSOR methods. **BETAB** may change depending on the adaptive switch **IPARM(6)**. The matrix L is the strictly lower triangular part of the Jacobi matrix and U the strictly upper triangular part. When the spectral radius of LU is less than or equal to $\frac{1}{2}$, the "SSOR condition" is satisfied for some problems, provided one uses the natural ordering. (See [4, 5, 19] for additional details.)
- RPARM(8)**
100. * SRELPR **TOL:** Tolerance factor near machine relative precision, **SRELPR**. In each row, if all nonzero off-diagonal row entries are less than **TOL** times the value of the diagonal entry, then this row and corresponding column are essentially removed from the system. This is done by setting the nonzero off-diagonal elements in the row and corresponding column to zero, replacing the diagonal element with 1, and adjusting the elements on the right-hand side of the system so that the new system is equivalent to the original one.⁵ If the diagonal entry is the only nonzero element in a row and is not greater than the reciprocal of **TOL**, then no elimination is done. This procedure is useful for linear systems arising from finite-element discretizations of PDEs in which Dirichlet boundary conditions are handled by giving the diagonal values in the linear system extremely large values. (The installer of this package should set the value of **SRELPR**. See comments in subroutine **DFAULT()** for additional details.)
- RPARM(9) 0.** **TIME1:** Total time in seconds from beginning of iterative algorithm until convergence. (A machine-dependent sub-program call for returning the time in seconds is provided by the installer of this package.)
- RPARM(10) 0.** **TIME2:** Total time in seconds for entire call.
- RPARM(11) 0.** **DIGIT1:** Approximate number of digits using the estimated relative error with the final approximate solution. Computed as the negative of the logarithm base 10 of the final value of the stopping test. (See details below or [6].)
- RPARM(12) 0.** **DIGIT2:** Approximate number of digits using the estimated relative residual with the final approximate solution. Computed as the negative of the logarithm base 10 of the ratio of the two norm of the residual vector and the two norm of the right-hand-side vector. This estimate is related to the condition number of the original linear system and therefore it will not be accurate if the system is ill-conditioned. (See details below or [6].)

⁵ If the row and column corresponding to diagonal entry $\text{coef}(i, i)$ are to be eliminated, then the right-hand side is adjusted to $\text{rhs}(i) = \text{rhs}(i)/\text{coef}(i, i)$ and $\text{rhs}(j) = \text{rhs}(j) - \text{rhs}(i) \times \text{coef}(i, j)$ for $j \neq i$.

DIGIT1 is determined from the actual stopping test computed on the final iteration, whereas **DIGIT2** is based on the computed residual vector using the final approximate solution after the algorithm has converged. If these values differ greatly, then either the stopping test has not worked successfully or the original system is ill-conditioned. (See [6] for additional details.)

For storage of certain intermediate results, the solution modules require a real vector **WKSP(*)** and a corresponding variable **NW** indicating the available space. The length of the work-space array varies with each solution module and the maximum amount needed is given in the following table.

<u>SOLUTION MODULE</u>	<u>MAXIMUM LENGTH OF WKSP(*)</u>
JCG()	$4 * N + NCG$
JSI()	$2 * N$
SOR()	N
SSORCG()	$6 * N + NCG$
SSORSI()	$5 * N$
RSCG()	$N + 3 * NB + NCG$
RSSI()	$N + NB$

where value of **NCG** is $2 * IPARM(1)$ for symmetric sparse storage and $4 * IPARM(1)$ for nonsymmetric sparse storage. It should be noted that the actual amount of work space used may be somewhat less than these upper limits since some of the latter are dependent on the maximum number of iterations allowed, **ITMAX**, stored in **IPARM(1)**. Clearly, the array **WKSP(*)** must be dimensioned to at least the value of **NW**.

Nonzero integer values of the error flag **IER** indicate that an error condition was detected. These values are listed below according to their numerical value and to the name of the routine in which the flag was set.

<u>ERROR FLAG</u>	<u>MEANING</u>
IER = 0,	Normal convergence obtained.
= 1 + Mth,	Invalid order of system, N
= 2 + Mth,	Work-space array WKSP(*) not large enough— IPARM(8) set to amount of required work space, NW .
= 3 + Mth,	Failure to converge in IPARM(1) iterations— RPARM(1) reset to last stopping value computed.
= 4 + Mth,	Invalid order of black subsystem, NB .
= 101,	Diagonal element not positive.
= 102,	No diagonal entry in a row.
= 201,	Red-black indexing not possible.
= 301,	No entry in row of original matrix.
= 302,	No entry in row of permuted matrix.
= 303,	Sorting error in row of permuted matrix.
= 401,	Diagonal element not positive.
= 402,	No diagonal entry in a row.
= 501,	Failure to converge in ITMAX evaluations.
= 502,	Function does not change sign at endpoints.
= 601,	Successive iterates not monotone increasing.

JCG(), **JSI()**, **SOR()**, **SSORCG()**, **SSORSI()**, **RSCG()**, **RSSI()** assign values to **Mth** of 10, 20, 30, 40, 50, 60, 70, respectively. **SBELM()**, **PRBNDX()**, **PERMAT()**, **SCAL()**, **ZBRENT()**, **EQRT1S()** are subroutines with error flags in the 100s, 200s, 300s, 400s, 500s, 600s, respectively. These routines perform the following functions: **SBELM()** removes rows and columns, **PRBNDX()** determines the red-black indexing, **SCAL()** scales the system, **ZBRENT()** is a modified IMSL routine for computing a zero of a function which changes sign in

a given interval, **EQRT1S()** is a modified IMSL routine for computing the largest eigenvalue of a symmetric tridiagonal matrix.⁶

4. USER-ORIENTED MODULES

The array **U(*)** should contain an initial approximation to the solution of the linear system before any **ITPACK** module is called. If the user has no information for making such a guess, then the zero vector may be used as the starting vector. The subroutine **VFILL()** can be used to fill a vector with a constant:

CALL VFILL (N, U, VAL)

fills array **U(*)** of length **N** with value **VAL** in each entry.

To aid the user in using the iterative methods of **ITPACK**, four modules for constructing the sparse-matrix storage arrays are included. The modules are

- SBINI()** called at the beginning to initialize the arrays **IA(*)**, **JA(*)**, **A(*)**, and **IWORK(*)**;
- SBSIJ()** called repeatedly to set the individual entries in the matrix and build a link-list representation of the matrix structure;
- SBEND()** called at the end to restructure the link list into final sparse storage form;
- SBAGN()** called to return again to the link-list representation if **SBEND()** has been called but additional elements are to be added or modified.

These modules are described below.

(a) Initialization:

CALL SBINI (N, NZ, IA, JA, A, IWORK)

Initializes **IA(*)**, **JA(*)**, **A(*)**, and **IWORK(*)** for a system of order **N**. **IA(*)**, **JA(*)**, and **IWORK(*)** are integer arrays of length at least **N + 1**, **NZ**, and **NZ**, respectively. **A(*)** is a real array of length at least **NZ**.

(b) Set individual entries:

CALL SBSIJ (N, NZ, IA, JA, A, IWORK, I, J, VAL, MODE, LEVEL, NOUT, IER)

Inserts the value, **VAL**, of the **(I, J)** entry of the user's matrix into the link-list representation for that matrix. When using symmetric sparse storage, **J** must be greater than or equal to **I**. If the **(I, J)** entry has already been set, then **MODE** specifies the way in which the entry is to be treated:

- MODE** < 0, Current entry value is left as is;
- = 0, Current entry value is reset to **VAL**;
- > 0, **VAL** is added to the current entry value.

If **LEVEL** is less than 0, **SBSIJ()** causes no printing. If **LEVEL** is 0, fatal error messages are written to output unit number **NOUT**; and if **LEVEL** is 1 or greater, a message is printed when **SBSIJ()** encounters a value it has already set with the value being reset according to the value of **MODE**. **IER** is an error parameter and returns values of

<u>ERROR FLAG</u>	<u>MEANING</u>
IER = 0,	New (I, J) entry established.
= 700,	(I, J) entry already set—reset according to MODE .
= 701,	Improper values for either I or J .
= 702,	NZ too small—no room for new entry.

⁶ International Mathematical and Statistical Libraries, Inc., Sixth Floor NBC Bldg., 7500 Bellaire Blvd., Houston, TX 77036.

(c) Finalization:

CALL SBEND (N, NZ, IA, JA, A, IWORK)

Restructures the link-list data structure built by **SBINI()** and **SBSIJ()** into the final data structure required by **ITPACK**.

(d) Undo finalization:

CALL SBAGN (N, NZ, IA, JA, A, IWORK, LEVEL, NOUT, IER)

Returns to link-list representation for modification or addition of elements to the system. Repeated calls to **SBSIJ()** can then be made followed by a single call to **SBEND()** to close out the sparse-matrix representation. If **LEVEL** is less than 0, no printing is done, and if **LEVEL** is 0 or greater, fatal error information is written to output unit number **NOUT**. **IER** is an error flag indicating

<u>ERROR FLAG</u>	<u>MEANING</u>
IER = 0,	Successful completion.
= 703,	NZ too small—no room for new entry.

Note that **SBINI()** should not be called after **SBAGN()** is called, since it would destroy the previous data.

5. EXAMPLES

Given a linear system $Au = b$ with

$$A = \begin{bmatrix} 4 & -1 & -1 & 0 \\ -1 & 4 & 0 & -1 \\ -1 & 0 & 4 & -1 \\ 0 & -1 & -1 & 4 \end{bmatrix}, \quad b = \begin{bmatrix} 6 \\ 0 \\ 0 \\ 6 \end{bmatrix}$$

a program to solve this problem with an initial guess of $u^T = (0, 0, 0, 0)$ using **JCG()** with symmetric sparse storage and printing the final approximate solution vector follows.

```

INTEGER IA(5), JA(8), IPARM(12), IWKSP(12)
REAL A(8), RHS(4), U(4), WKSP(24), RPARAM(12)
DATA A(1), A(2), A(3), A(4)/4.0, -1.0, -1.0, 4.0/
DATA A(5), A(6), A(7), A(8)/-1.0, 4.0, -1.0, 4.0/
DATA JA(1), JA(2), JA(3), JA(4)/1, 2, 3, 2/
DATA JA(5), JA(6), JA(7), JA(8)/4, 3, 4, 4/
DATA IA(1), IA(2), IA(3), IA(4), IA(5)/1, 4, 6, 8, 9/
DATA RHS(1), RHS(2), RHS(3), RHS(4)/6.0, 0.0, 0.0, 6.0/
DATA N/4/, NW/24/, ITMAX/4/, LEVEL/1/, IDGTS/2/

```

C

```

CALL DFAULT (IPARM, RPARAM)
IPARM(1) = ITMAX
IPARM(2) = LEVEL
IPARM(12) = IDGTS
CALL VFILL (N, U, 0.E0)
CALL JCG (N, IA, JA, A, RHS, U, IWKSP, NW, WKSP, IPARM, RPARAM, IER)
STOP
END

```

The output for this run would be

```

BEGINNING OF ITPACK SOLUTION MODULE JCG
JCG HAS CONVERGED IN 2 ITERATIONS.
APPROX. NO. OF DIGITS (EST. REL. ERROR) = 14.6 (DIGIT1)
APPROX. NO. OF DIGITS (EST. REL. RESIDUAL) = 14.3 (DIGIT2)
SOLUTION VECTOR.

```

1	2	3	4
2.00000E+00	1.00000E+00	1.00000E+00	2.00000E+00

Textbook methods such as the Jacobi (J), Gauss-Seidel (GS), successive overrelaxation (SOR—fixed relaxation factor ω), symmetric successive

Table I

Method	Use	Parameters
J	JSI()	IPARM(6) = 0, IPARM(7) = 2
GS	SOR()	IPARM(6) = 0
SOR—fixed omega	SOR()	IPARM(6) = 0, RPARAM(5) = OMEGA
SSOR—fixed omega	SSORSI()	IPARM(6) = 0, RPARAM(5) = OMEGA
RS	RSSI()	IPARM(6) = 0

overrelaxation (SSOR—fixed relaxation factor omega), and the RS method can be obtained from this package by resetting appropriate parameters after the subroutine **DFAULT()** is called but before **ITPACK** routines are called (see Table I). These methods were not included as separate routines because they are usually slower than the accelerated methods included in this package.

On the black unknowns, the cyclic Chebyshev semi-iterative (CCSI) method of Golub and Varga [2] gives the same result as the RSSI method. The CCSI and RSSI methods converge at the same rate, and each of them converges twice as fast as the JSI method. This is a theoretical result [6] and does not count the time involved in establishing the red-black indexing and the red-black partitioned system. Similarly, the cyclic conjugate gradient (CCG) method with respect to the black unknowns, considered by Reid [16] (see also Hageman and Young [6]), gives the same results as the RSCG method. Also, the CCG and the RSCG methods converge at the same rate, and each of them converges, theoretically, exactly twice as fast as the JCG method. Hence, the accelerated RS methods are preferable to the accelerated J methods when using a red-black indexing.

6. NUMERICAL RESULTS

The iterative algorithms in **ITPACK** have been tested over a wide class of matrix problems arising from elliptic partial differential equations with Dirichlet, Neumann, and mixed boundary conditions on arbitrary two-dimensional regions (including cracks and holes) and on rectangular three-dimensional regions [1]. Both finite-difference and finite-element procedures have been employed to obtain the linear systems. The two sample problems presented here, while simple to pose, are representative of the behavior of the **ITPACK** routines for more complex problems. The iterative algorithms make no use of the constant coefficients in these two problems or of the particular structure of the resulting linear system. Because the **ITPACK** code is not tailored to any particular class of partial differential equations or discretization procedure, but rather to sparse linear systems, it is felt that the package can be used to solve a wider class of problems.

We now consider two simple partial differential equations that, when discretized by finite-difference methods, give rise to large sparse linear systems. We obtain the solution of each of these systems by the seven algorithms in **ITPACK 2C**. These numerical results should aid the user of **ITPACK** in determining the amount of time required when solving more complicated sparse systems. However, one should not interpret these execution times as conclusive by themselves. Variances introduced by different compilers, computer systems, and timing functions can sometimes be significant. Moreover, the number of iterations required by an iterative method is dependent on the problem being solved, the initial estimate for the solution, the parameter estimates used, and the relative accuracy requested in the stopping criterion **RPARAM(1)**. These tests were run on the CDC CYBER 170/750 at the University of Texas with the FTN 4.8 compiler (**OPT** = 2).

To obtain representative sparse linear systems, we discretize the following two self-adjoint elliptic partial differential equations in a region with prescribed conditions on the boundary. Here u_{xx} , u_{yy} , u_{zz} are partial derivatives and du/dn is the derivative in the normal direction.

Table II. Number of Iterations and Execution Times for Problem (1) Using Adaptive and Nonadaptive Procedures (Nonadaptive Data in Parentheses)

Routine	Ordering	Iterations	Iteration time	Total time
JCG()	Natural	61 (61)	0.250 (0.247)	0.281 (0.271)
	Red-black	61 (61)	0.232 (0.246)	0.402 (0.413)
JSI()	Natural	108 (95)	0.408 (0.344)	0.439 (0.375)
	Red-black	108 (95)	0.393 (0.332)	0.569 (0.498)
SOR()	Natural	72 (54)	0.356 (0.280)	0.368 (0.307)
	Red-black	65 (47)	0.311 (0.224)	0.469 (0.411)
SSORCG()	Natural	17 (13)	0.232 (0.173)	0.264 (0.185)
SSORSI()	Natural	23 (22)	0.242 (0.213)	0.273 (0.244)
RSCG()	Red-black	31 (31)	0.104 (0.117)	0.269 (0.297)
RSSI()	Red-black	60 (48)	0.207 (0.166)	0.358 (0.344)

Table III. Number of Iterations and Execution Times for Problem (2) Using Adaptive and Nonadaptive Procedures (Nonadaptive Data in Parentheses)

Routine	Ordering	Iterations	Iteration time	Total time
JCG()	Natural	28 (28)	0.092 (0.090)	0.107 (0.090)
	Red-black	28 (28)	0.079 (0.074)	0.191 (0.202)
JSI()	Natural	64 (54)	0.166 (0.136)	0.196 (0.152)
	Red-black	64 (54)	0.160 (0.130)	0.268 (0.266)
SOR()	Natural	42 (29)	0.139 (0.095)	0.150 (0.110)
	Red-black	38 (29)	0.124 (0.097)	0.236 (0.231)
SSORCG()	Natural	15 (11)	0.136 (0.097)	0.167 (0.111)
SSORSI()	Natural	19 (15)	0.138 (0.101)	0.153 (0.117)
RSCG()	Red-black	15 (15)	0.032 (0.051)	0.150 (0.169)
RSSI()	Red-black	31 (27)	0.075 (0.064)	0.186 (0.196)

$$\begin{cases} u_{xx} + 2u_{yy} = 0, & (x, y) \text{ in } S = (0, 1) \times (0, 1) \\ u = 1 + xy, & (x, y) \text{ on boundary of } S. \end{cases} \quad (1)$$

Using the standard 5-point symmetric finite-difference operator with $h = \frac{1}{20}$, we obtain a sparse linear system with 1729 nonzero elements and 361 unknowns.

$$\begin{cases} u_{xx} + 2u_{yy} + 3u_{zz} = 0, & (x, y, z) \text{ in } C = (0, 1) \times (0, 1) \times (0, 1) \\ \text{On boundary of } C: \\ \quad u = 1, & (0, y, z), (x, 0, z), \text{ or } (x, y, 0) \\ \quad \left\{ \frac{du}{dn} = \begin{cases} yz(1 + yz), & (1, y, z) \\ xz(1 + xz), & (x, 1, z) \\ xy(1 + xy), & (x, y, 1). \end{cases} \right. \end{cases} \quad (2)$$

Using the standard 7-point symmetric finite difference operator with $h = \frac{1}{7}$, we obtain a sparse linear system with 1296 nonzero elements and 216 unknowns.

Tables II and III display the number of iterations and execution times (in seconds) for the seven methods in ITPACK 2C for the linear systems corresponding to problems (1) and (2), respectively, using symmetric sparse storage. Both the time for the iteration algorithm and the total time for the subroutine call are given. The stopping criterion was set to $5.0E-6$. To illustrate how effective the adaptive procedures are, we have included in these tables the number of iterations and the time when the optimum iteration parameters were used with no adaptive procedures.

Values corresponding to the red-black ordering with the SSOR methods are omitted from the tables since it is known that these methods are ineffective with

this ordering. Since the RS methods are defined for only the red-black ordering, the table entries for these methods with the natural ordering are not included.

7. NOTES ON USE

Before an iterative algorithm is called to solve a linear system, the values in the array **A(*)** are permuted and scaled. Afterward, these values are unpermuted and unscaled. Consequently, the values in arrays **A(*)** and **RHS(*)** may change slightly due to round-off errors in the computer arithmetic. Moreover, since entries in each row of the linear system may be stored in any order within a contiguous block of data, the locations of elements of **A(*)** and of corresponding ones in **JA(*)** may change from those given before the permuting and unpermuting was done. The same linear system is defined by the arrays **A(*)**, **JA(*)**, and **IA(*)**, whether or not corresponding elements in **A(*)** and **JA(*)** have changed locations within contiguous blocks.

Scaling of the linear system is done as follows to reduce the number of arithmetic operations. The diagonal entries of the linear system are checked for positivity and are moved to the first **N** locations of the array **A(*)**. The nonzero off-diagonal entries of the linear system $Au = b$ are scaled. The scaling involves the diagonal matrix $D^{1/2}$ of square roots of the diagonal entries of the linear system, that is,

$$(D^{-1/2}AD^{-1/2})(D^{1/2}u) = (D^{-1/2}b).$$

The algorithms iterate until convergence is reached based on the relative accuracy requested via the stopping criterion set in **RPARM(1)** for the scaled solution vector $(D^{1/2}u)$. Unscaling solves for u and returns the linear system to its original form subject to round-off errors in the arithmetic and to possible movement of entries within contiguous blocks of data.

When requested, a red-black permutation of the data will be done before and after the iterative algorithm is called. Otherwise, the linear system is used in the order it is given, which we call the "natural ordering."

The successive overrelaxation (SOR) method has been shown to be more effective with the red-black ordering than with the natural ordering for some problems [19]. In the SOR algorithm, the first iteration uses **OMEGA = 1.0** and the stopping criterion is set to a large value so that at least one Gauss-Seidel iteration is performed before an approximate value for the optimum relaxation parameter is computed.

Optional features of this package are red-black ordering, effective removal of rows and columns when the diagonal entry is extremely large, and error analysis. In the event that one is not using some of these options and needs additional memory space for a very large linear system, the relevant subroutines that can be replaced with dummy subroutines are as follows: red-black ordering [**PRBNDX()**, **PERMAT()**, **PERVEC()**, **QSORT()**], removal of rows [**SBELM()**], error analysis [**PERROR()**].

The timing routines **ITICK()** and **ITOCK()** should call a local system routine so that they return the run time in milliseconds. **ITICK()** is called at the beginning of a timing interval and **ITOCK()** at the end.

The value of the machine relative precision is contained in the variable **SRELPR**, which is set in the subroutine **DFAULT()** and in the test program. This and other default values may be permanently changed when the code is installed by changing their values in the subroutine **DFAULT()**. **SRELPR** must be changed when moving the code to another computer. If the installer of this package does not know its value, an approximate value can be determined from a simple FORTRAN program given in the comment statements of subroutine **DFAULT()**.

Since the amount of precision may change from computer to computer, the relative accuracy requested in the stopping criterion **ZETA** must not be less than about **500**. times the machine relative precision **SRELPR**. If a value of **ZETA** is

requested that is too small, then the code resets it to this value. The current default value for **ZETA**, **5.0E-6**, is set by the routine **DFAULT()** into **RPARM(1)**.

The distribution tape contains the ITPACK 2C software package of 71 subprograms and a testing program **MAIN()** together with its 27 subprograms. The routines **DFAULT()**, **ITICK()**, **ITOCK()** in ITPACK and the program **MAIN()** are the only ones requiring editing by the installer of the package. ITPACK can be made into a compiled program library, although not all of it would normally be used in a particular application.

8. ITPACK HISTORY

The 2C version of the ITPACK codes described here is the result of several years of research and development. The development of ITPACK began in the early 1970s when Professor Garrett Birkhoff suggested that general-purpose software for solving linear systems should be developed for iterative methods as well as for direct methods. Initially, prototype programs were written based on preliminary iterative algorithms involving adaptive selection of parameters and automatic stopping procedures. These programs were tested on a large set of elliptic partial differential equations over domains compatible with the subroutine **REGION()** [8], which superimposed a square grid over the domain. These routines were designed for solving self-adjoint elliptic partial differential equations. Next a preliminary version of ITPACK was coded in standard FORTRAN. The ITPACK routines used iterative algorithms that were refined from the prototype programs. However, these routines were designed to solve large sparse linear systems of algebraic equations instead of partial differential equations. The use of three interchangeable symmetric sparse storage modes in ITPACK 1.0 [3] allowed for great flexibility and made it possible to solve a wider class of problems than the prototype programs and to study different storage modes for iterative methods. The next version, ITPACK 2.0 [4], was significantly faster than its predecessor, since it was restricted to allow only one sparse symmetric storage format. Most of the iterative algorithms utilized in the 2.0 version of this package assume that the coefficient matrix of the linear system is symmetric positive definite. As with many packages, the need to handle a slightly larger class of problems, namely, nearly symmetric systems, soon became evident. This required adapting the routines to allow a switch for either a symmetric or nonsymmetric storage mode in ITPACK 2A [5]. Moreover, a modification of the conjugate gradient algorithms was developed to handle nearly symmetric systems [13]. ITPACK has been improved in the 2B version [11] by (a) writing more efficient versions of several key subroutines, (b) incorporating basic linear algebra subprograms (BLAS) [15], and (c) improving the user interface with better printing and documentation. Some additional improvements and corrections were made in the 2C version. The algorithms in ITPACK are not guaranteed to converge for all linear systems, but have been shown to work successfully for a large number of symmetric and nonsymmetric systems that arise from solving elliptic partial differential equations [1, 14].

The numerical algorithms in ITPACK 2C correspond to those described in the appendix of technical report [4] and outlined in the book [6]. In particular, the SOR code is based on an algorithm suggested to us by L. Hageman. Various other algorithms exist for iterative methods. For example, S. Eisenstat has an implementation of the symmetric successive overrelaxation preconditioned conjugate gradient procedure.⁷

Modules based on the seven iterative routines in ITPACK have been incorporated into the elliptic partial differential equation solving package ELLPACK [17] together with all the necessary translation routines needed. The user-oriented modules described in Section 4 are not in ELLPACK. Moreover, if the ELLPACK

⁷ Private communication.

system is not being used to generate the linear system for ITPACK, it is recommended that ITPACK be used as a stand-alone package apart from ELLPACK.

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ALGORITHM 587

Two Algorithms for the Linearly Constrained Least Squares Problem

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Categories and Subject Descriptors: G.1.2. [Numerical Analysis]: Approximation—*least squares approximation*; G.1.6. [Numerical Analysis]: Optimization—*constrained optimization*; *least squares methods*; G.m [Mathematics of Computing]: Miscellaneous—*FORTRAN*

General Terms: Algorithms

Additional Key Words and Phrases: linear least squares solution, equality constraints, inequality constraints, nonnegativity constraints, inconsistent constraints, covariance matrix

1. INTRODUCTION

This paper discusses subroutines for computing numerical solutions of the following two linearly constrained linear least squares problems.

$$\begin{array}{ll}
 \textit{Problem NNLSE} & \begin{array}{l} E\mathbf{x} = \mathbf{f} \quad (\text{equations to be exactly satisfied}) \\ A\mathbf{x} \cong \mathbf{b} \quad (\text{equations to be approximately} \\ \quad \quad \quad \text{satisfied, least squares sense}) \\ x_i \geq 0, \quad i = l + 1, \dots, n, \quad 0 \leq l \leq n \end{array}
 \end{array} \tag{1}$$

$$\begin{array}{ll}
 \textit{Problem LSEI} & \begin{array}{l} E\mathbf{x} = \mathbf{f} \quad (\text{equations to be exactly satisfied}) \\ A\mathbf{x} \cong \mathbf{b} \quad (\text{equations to be approximately} \\ \quad \quad \quad \text{satisfied, least squares sense}) \\ G\mathbf{x} \geq \mathbf{h} \quad (\text{inequality constraints that the} \\ \quad \quad \quad \text{solution must satisfy}) \end{array}
 \end{array} \tag{2}$$

In both problems the matrices E and A are real and of respective dimensions m_E by n and m_A by n . For Problem NNLSE, the variables x_1, \dots, x_l are free to have either sign. For Problem LSEI, the (real) inequality constraint matrix G is m_G by n . The right-side vectors \mathbf{f} , \mathbf{b} , and \mathbf{h} that appear in the two problem statements have, respectively, m_E , m_A , and m_G components. The (unknown) solution vector \mathbf{x} has n components.

While Problem LSEI of eq. (2) appears to be a more general problem than Problem NNLSE of eq. (1), it really is not. In fact, there are a number of ways to transform Problem LSEI into one of the forms of Problem NNLSE. Three ways of doing this are discussed in [3]. The method we have implemented is described on pages 101–102. The successful implementation of an algorithm for solving

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Problem NNLSE is the key computational process. Nevertheless, it is important for applications such as constrained curve fitting [2] to have a subprogram that solves Problem LSEI of eq. (2) directly. We provide FORTRAN subprograms **WNNLS()** and **LSEI()** that solve the respective problems in eqs. (1) and (2).

In Section 2 we review mathematical and numerical analysis details pertinent to solving Problem LSEI. In Section 3 we review some necessary details for understanding our methods for solving Problem NNLSE. In Section 4 we summarize some features and advantages of the codes. These features include changing tolerances, scaling of data matrices, and optional computation of the covariance matrix. Section 5 presents a test subprogram **CLSTP()**, which is included with the package. It solves the test problem with both subprograms. Section 6 contains installation guidelines and remarks.

2. SOLVING PROBLEM LSEI

In this section, we briefly review mathematical and algorithmic details needed to solve Problem LSEI of eq. (2) [3, pp. 101-102]. The overall process consists of four main parts.

- Step 1 Problem LSEI is reduced to a subproblem with possibly fewer unknown variables and with all explicitly stated equality constraints removed.
- Step 2 The problem resulting from step 1 is reduced to a new problem where the least squares matrix is a simple projection matrix and the right-side vector is zero.
- Step 3 The problem resulting from step 2 is solved by reposing it as a dual problem. This dual problem consists of two special cases of Problem NNLSE, eq. (1).
- Step 4 The solution obtained in step 3 is transformed to the solution of the original problem using translations, matrix multiplications, and the solution of triangular linear algebraic systems.

3. SOLVING PROBLEM NNLSE

The theoretical development for solving problem NNLSE of eq. (1) is presented in [3]. The fundamental point of this method involves a numerically stable implementation of a penalty function approach. The least squares equations are each weighted by a small parameter ϵ , chosen in the subprogram **WNNLS()**. The augmented and weighted least squares system of eq. (3) is then solved.

$$\begin{bmatrix} E \\ \epsilon A \end{bmatrix} \mathbf{x} \equiv D \begin{bmatrix} E \\ A \end{bmatrix} \mathbf{x} \cong D \begin{bmatrix} \mathbf{f} \\ \mathbf{b} \end{bmatrix} \equiv \begin{bmatrix} \mathbf{f} \\ \epsilon \mathbf{b} \end{bmatrix}$$

$$D = \text{diag} \left(\overbrace{1, \dots, 1}^{m_E}, \overbrace{\epsilon, \dots, \epsilon}^{m_A} \right) \quad (3)$$

$$\mathbf{x} = \begin{bmatrix} \mathbf{y} \\ \mathbf{w} \end{bmatrix} \begin{matrix} \} l \\ \} n - l \end{matrix} \quad \begin{matrix} \mathbf{y} \text{ unconstrained} \\ \mathbf{w} \geq \mathbf{0} \end{matrix}$$

Part of the theoretical development in [3] shows that solutions of the weighted problem of eq. (3) converge to solutions of Problem NNLSE (if it is consistent) as $\epsilon \rightarrow 0$. Within the subprogram **WNNLS()** eq. (3) is solved only once with a value of ϵ that is chosen to achieve full working accuracy in the solution. The value used in **WNNLS()** is defined by

$$\epsilon^2 = \frac{10^{-4}\eta}{\gamma}, \quad (4)$$

where $\gamma = \| \frac{E}{A} \|$, ($\| \cdot \|$ = subordinate matrix norm of l_∞ vector norm), and η = machine relative arithmetic precision.

The algorithm for solving eq. (3) with ϵ as defined in eq. (4) proceeds in two main steps. First we compute a (minimum-length) solution for the unconstrained

variables in terms of the constrained variables. Solving for the unconstrained variables is primarily a triangularization operation. In the second main step of the process we solve for the constrained variables. This is an iterative process, that is, it is Algorithm NNLS of [7, Chap. 23]. Certain crucial differences in numerical tests are needed because of the penalty parameter ϵ that multiplies the least squares equations. These tests are discussed in [3].

4. USAGE SUGGESTIONS AND SUBPROGRAM OPTIONS

In Sections 2 and 3 we have outlined solution methods for solving Problem LSEI of eq. (2) and Problem NNLS of eq. (1). As shown in [3], computing the solution of Problem NNLS can be regarded as the core computation in solving constrained linear least squares problems.

The most satisfactory method from the standpoint of accuracy and stability is to introduce slack variables into the inequality constraints of Problem LSEI [3]. This problem is then solved using subprogram WNNLS(). The results of solving a bounded variable Hilbert matrix problem summarized in [3] suggest that subprogram WNNLS() continues to compute acceptable solutions even as the problems become increasingly ill-conditioned.

The use of subprogram WNNLS() with the slack variable formulation does have a disadvantage compared to subprogram LSEI(). For most problems, WNNLS() will require more computing time and storage than LSEI(). This is due to the larger number of problem variables in the slack variable formulation. The advantage of efficiency with LSEI() may be countered by the simultaneous occurrence of poor conditioning and rounding errors. (This can occur with a poorly conditioned least squares problem.) Owing to the poor conditioning and rounding error, the feasible constraint region can be mapped to one that is infeasible. Instances of this are shown in the results of solving the bounded variable Hilbert matrix problem summarized in [3].

The choice between the two subprograms is a time and storage versus stability trade-off. Specifically, in the case of a poorly conditioned least squares problem, WNNLS() might obtain a solution when LSEI() cannot. As illustrated in [3], subprogram WNNLS() can also be used to extend the notion of solution for problems with infeasible constraints.

Occasionally, a user of subprogram LSEI() will need the covariance matrix of the least squares solution variables of minimum length. This is returned as an output matrix if the user wants it. It is an unbiased estimate of the covariance matrix for the minimum-length solution of an equality constrained least squares problem *with no inequalities*. This is developed in [4] and [6].

When inequalities are included, certain additional mathematical problems must be considered. These have to do with the behavior of the set of inequalities chosen by the algorithm to be equalities. The question is as follows: What is the sensitivity of these equalities as the data are allowed to vary within its uncertainty? Inequalities may move from being satisfied as equalities to strict inequalities as the data are perturbed. The covariance matrix computed by LSEI() is based on the assumption that the set of equalities does not change when the solution is perturbed. No comprehensive theory is known to the authors for determining the matrix when the set of equalities does change. The user must keep these facts in mind when interpreting the covariance matrix for Problem LSEI with inequalities.

The remainder of this section describes parameters within LSEI() and WNNLS() which can optionally be changed by the user. These options fall into the three following groups.

- (A) Computation of the covariance matrix.
- (B) Column scaling of the data matrix.
- (C) Redefinition of tolerances used for determining ranks of problem matrices.

Changes to any number of these parameters can be specified as the linked-list input in the array **PRGOPT**(*). Precise instructions for defining **PRGOPT**(*) are found in the usage prologues for **LSEI**() and **WNNLS**(). If the user is satisfied with the nominal subprogram features, it is only necessary to set **PRGOPT**(1)=1.

Remarks about A: Nominally the covariance matrix is not computed by **LSEI**().

Remarks about B: Column scaling of the form $\mathbf{x} := D\mathbf{y}$ is always performed by **LSEI**() and **WNNLS**(). Nominally D is the identity matrix. Another option here is a choice for D such that each nonzero column of the entire scaled data matrix has length one. The user can also specify an arbitrary D .

Remarks about C: The user can change tolerances t_E and t_A in **LSEI**() and tolerance t_W in **WNNLS**(). The nominal values of t_E , t_A , and t_W are $\eta^{1/2}$, where η is the relative arithmetic precision of the machine. The parameter t_E is used in approximating the rank of the equality constraint matrix E of eq. (2). Its role is discussed near the end of [3, Sec. 1].

The parameter t_A is used in approximating the rank of the least squares matrix that results from eliminating the equality constraints from eq. (2). It is used to compute the factor τ , which is t_A times the norm of this reduced least squares matrix. Then τ is used in Algorithm HFTI [7, Chap. 14].

The parameter t_W is used by **WNNLS**() to compute the rank of the row-scaled least squares matrix as discussed in [3, Sec. 3.1].

5. REMARKS ON THE TESTING SUBPROGRAM **CLSTP**()

The subprogram **CLSTP** (**KLOG**, **COND**, **ISTAT**) constructs and solves a constrained least squares problem that has a known solution and known condition numbers [7, Chap. 9]. The problem generated is stated in eq. (2). The matrices A , E , and G are computed using formulas

$$A = U_1 S_1 V_1^T$$

$$E = U_2 S_2 V_2^T$$

and

$$G = U_3 S_3 V_3^T.$$

The problem dimensions are specified by using five integer parameters k_A , k_E , k_G , k_I , and k_n to compute $m_A = 2^{k_A}$, $m_E = 2^{k_E}$, $m_G = 2^{k_G}$, and $n = 2^{k_n}$. The integer $m_I = 2^{k_I}$ denotes the number of inequality constraints that are to be satisfied as strict inequalities. These five integers are passed to **CLSTP**() in the array **KLOG**(*) in the order indicated. If any of the values k_A , k_E , k_G , k_I , or k_n are less than zero, the respective values m_A , m_E , m_G , m_I , or m_n are set to zero. No computation is performed if $n = 0$.

Arrays within **CLSTP**() currently have fixed dimensions that require k_A , k_E , k_G , k_I , and k_n to all be less than or equal to 5. Instructions for increasing the array dimensions are given as comments within **CLSTP**().

The matrices U_j and V_j are symmetric orthogonal Hadamard matrices of dimension $n = 2^k$ generated by the recursion

$$n := 1$$

$$U := 1$$

For $l = 1, \dots, k$

$$U := \begin{bmatrix} U & : & U \\ U & : & -U \end{bmatrix}$$

$$n := n + n$$

End For

$$U := n^{-1/2}U$$

The matrices S_J , $J = 1, 2, 3$, are rectangular diagonal matrices. The extreme diagonal terms are κ_i and 1, where $\kappa_J \equiv \text{COND}(\mathbf{J})$. The intermediate diagonal terms are generated in the open interval $(1, \kappa_J)$ using the random number generator **RAN**(). The output value of $t = \text{RAN}(\text{ISEED})$ satisfies $0 < t < 1$. The intermediate diagonal terms are successively computed as $1 + t(\kappa_J - 1)$. Initially, **ISEED** is set to 100001 in **CLSTP**().

The n -vector $\hat{\mathbf{x}} = (1, \dots, 1)^T$ is used to generate the vectors

$$\mathbf{f} = E\hat{\mathbf{x}}$$

$$\hat{\mathbf{b}} = A\hat{\mathbf{x}}$$

and

$$\hat{\mathbf{h}} = G\hat{\mathbf{x}}.$$

We add a vector $\tilde{\mathbf{b}}$ to $\hat{\mathbf{b}}$ that is orthogonal to the column space of A . This is given by

$$\tilde{\mathbf{b}} = U_1(0, \dots, 0, g_{n+1}, \dots, g_{m_A}),$$

where

$$g_i = \text{RAN}(\text{ISEED}) \cdot \|\hat{\mathbf{b}}\| \cdot \sigma, \quad i = n + 1, \dots, m_A.$$

The value of σ is specified by the variable **ANSR** in **CLSTP**(). It is currently set to 0.01. The right-side vector for the least squares equations in eq. (2) is $\mathbf{b} \equiv \hat{\mathbf{b}} + \tilde{\mathbf{b}}$.

The right-side vector for the inequality constraints is constructed by making the first m_I constraints strict inequalities. This is done by defining the right-side vector as $\mathbf{h} = \hat{\mathbf{h}} - \tilde{\mathbf{h}}$, where

$$\tilde{\mathbf{h}} = (h_1, \dots, h_{m_I}, 0, \dots, 0^T),$$

$$h_i = \text{RAN}(\text{ISEED}) \cdot \|\hat{\mathbf{h}}\|, \quad i = 1, \dots, m_I.$$

These techniques for generating problems with known solutions are similar to those discussed in [9, pp. 6–9]. One might obtain different sets of test problems on machines with differing arithmetic characteristics. Part of this is due to a different sequence of numbers generated by **RAN**().

We have found that column scaling is sometimes required for solving eqs. (1) and (2). In particular, when using 32-bit floating-point arithmetic, problems generated by **CLSTP**() using the published test data occasionally failed to pass the tests when no column scaling was done. Thus the option array input for calls to both **LSEI**() and **WNNLS**() are set so that unit length column scaling is performed on all the tests.

After subprogram **LSEI**() has computed an approximate solution \mathbf{x}' for this particular form of eq. (2), and subprogram **WNNLS**() has solved for an approximate solution \mathbf{x}'' of the system

$$E\mathbf{x} = \mathbf{f}$$

$$A\mathbf{x} \cong \mathbf{b}$$

$$G\mathbf{x} - \mathbf{h} = \mathbf{w}$$

for the unknown $(\mathbf{x}^T, \mathbf{w}^T)^T$, we compute the differences $\mathbf{dx}_1 = \mathbf{x}' - \hat{\mathbf{x}}$ and $\mathbf{dx}_2 = \mathbf{x}'' - \hat{\mathbf{x}}$. A test is made on the value of $\|\mathbf{dx}_i\|$ to ensure that \mathbf{x}' or \mathbf{x}'' is as accurate as it deserves to be. The test of the subprogram has failed if the corresponding $\|\mathbf{dx}_i\|$ is too large. Otherwise the test has passed and \mathbf{x}' or \mathbf{x}'' is an acceptable approximation of \mathbf{x} . With

$$\rho = \|\tilde{\mathbf{b}}\| / \|\hat{\mathbf{b}}\|$$

$$\kappa = \kappa_1 = \text{condition number of } A$$

η = relative arithmetic precision

μ = $\max(m_A, n)$

ν = $\min(m_A, n)$

ϕ = 100

each test has passed if and only if

$$\frac{\|\mathbf{dx}_i\|}{\|\hat{\mathbf{x}}\|} \leq \kappa(1 + \kappa\rho)\eta[(6\mu - 3\nu)\nu]\phi.$$

The output value of **ISTAT** is set as follows:

ISTAT = 1 means both **LSEI**() and **WNNLS**() failed.
 = 2 means **WNNLS**() passed but **LSEI**() failed.
 = 3 means **LSEI**() passed but **WNNLS**() failed.
 = 4 means both **LSEI**() and **WNNLS**() passed.

This measure for $\|\mathbf{dx}_i\|$ is based on combining the estimate for the norm of the matrix H of the nearby problem that \mathbf{x}' solves (without constraints), $(A + H)\mathbf{x}' \cong \mathbf{b}$, [7, Chap. 13], together with the perturbation bounds of [7, Chap. 9].

It may be necessary to increase the value of ϕ slightly on some machines.

A short main program, **CLSTST**, is provided with the algorithm. Also provided are 11 data cards that are read by **CLSTST** from FORTRAN unit = 5. Each pair of the first 10 cards specifies a distinct test case. The last (eleventh) card terminates the program execution.

The subprogram **CLSTP**() prints the computed values of the least squares residual vector length and the vectors \mathbf{dx}_i for both **WNNLS**() and **LSEI**(). Also printed in **CLSTP**() are the computed ranks of the equality constraint and reduced least squares matrices returned by **LSEI**(). The arrays **KLOG(I)**, $I = 1$ to 5, and **COND(I)**, $J = 1$ to 3, and the value of **ISTAT** returned from **CLSTP**() are printed by **CLSTST**. Printing is done on FORTRAN unit = 6.

6. INSTALLATION GUIDELINES AND REMARKS

This section contains information for installing subprograms **LSEI**() and **WNNLS**().

Included in the package are seven groups of subprograms.

- (1) **LSEI, LSI, LPDP**
- (2) **WNNLS, WNLSM, WNLIT**
- (3) **HFTI, H12, DIFF** from [7]
- (4) **SDOT, SSCAL, SASUM, SAXPY, SNRM2, SCOPY, SSWAP, ISAMAX, SROTM, SROTMG** from [8]. (For double-precision usage **DDOT, DSCAL, DASUM, DAXPY, DNRM2, DCOPY, DSWAP, IDAMAX, DROTM, DROTMG**.)
- (5) **XERROR, XERRWV, XERABT, XERCLR, XERCTL, XERDMP, XERMAX, XERPRT, XERSAV, XGETF, XGETUA, XGETUN, XSETF, XSETUA, XSETUN, FDUMP, J4SAVE, S88FMT, NUMXER** from [5]. (The subprogram **NUMXER** is included for completeness but is not used in this package.)
- (6) **IIMACH** based on [1]
- (7) **CLSTST, CLSTP, RAN** (test package)

All of the subprograms are written in 1966 American National Standard portable FORTRAN. The only machine-sensitive subprogram is **IIMACH**(). It provides two environmental parameters required by the error-handling subprograms **XERROR**() and **XERRWV**(). This will require modification of **IIMACH**() at each host site. FORTRAN **DATA** statements defining the values of all the required constants are available for many machines in comments within the subprogram. The appropriate set of commented statements must be activated. If the values for your machine are not there, they should be provided in the order

corresponding to the description near the beginning of **IIMACH**(). Machines for which these constants are provided are Honeywell 600/6000, IBM 360/370, Xerox Sigma, CDC 6000/7000, PDP-10 (KA and KI processors), PDP-11 (16- and 32-bit arithmetic), Burroughs 5700/6700/7700, UNIVAC 1100, Data General Eclipse, Harris, VAX, and CRAY. In addition, the user must open or declare the FORTRAN unit, designated in **IIMACH**(4), where any error messages will be written.

We strongly recommend that calls to the error-handling subprograms **XERROR**() and **XERRWV**() be left intact. If the size or complexity of the error-handling package presents a problem on a particular machine, we suggest that the subprograms **XERROR**() and **XERRWV**() be replaced by shorter, machine-sensitive versions. These replacements should, minimally, print the character string comprising the error message and the specified data values. Usage of the full error-handling package is discussed in [5].

To convert the package for double-precision usage, follow the editing instructions at the beginning of each subprogram in groups 1, 2, 3, and 7 above. Use the double-precision version of the BLAS in group 4. No conversion is required for subprograms in groups 5 and 6.

ACKNOWLEDGMENTS

We thank two anonymous referees for providing suggestions that clarified the presentation of this algorithm.

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ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```

SUBROUTINE LSEI(W, MDW, ME, MA, MG, N, PRGOPT, X, RNORME, RNORML, LSEI 10
* MODE, WS, IP)                                     LSEI 20
C                                                     LSEI 30
C DIMENSION W(MDW, N+1), PRGOPT(*), X(N),           LSEI 40
C WS(2*(ME+N)+K+(MG+2)*(N+7)), IP(MG+2*N+2)       LSEI 50
C ABOVE, K=MAX(MA+MG, N).                           LSEI 60
C                                                     LSEI 70
C ABSTRACT                                           LSEI 80
C                                                     LSEI 90
C THIS SUBPROGRAM SOLVES A LINEARLY CONSTRAINED LEAST SQUARES LSEI 100
C PROBLEM WITH BOTH EQUALITY AND INEQUALITY CONSTRAINTS, AND, IF THE LSEI 110

```


C	(WT*E)X = (WT*F)	WNN	250
C	(A) (B), (LEAST SQUARES)	WNN	260
C	SUBJECT TO COMPONENTS L+1,...,N NONNEGATIVE.	WNN	270
C		WNN	280
C	THE SUBPROGRAM CHOOSES THE HEAVY WEIGHT (OR PENALTY PARAMETER) WT.	WNN	290
C		WNN	300
C	THE PARAMETERS FOR WNNLS ARE	WNN	310
C		WNN	320
C	INPUT..	WNN	330
C		WNN	340
C	W(*,*),MDW, THE ARRAY W(*,*) IS DOUBLE SUBSCRIPTED WITH FIRST	WNN	350
C	ME,MA,N,L DIMENSIONING PARAMETER EQUAL TO MDW. FOR THIS	WNN	360
C	DISCUSSION LET US CALL M = ME + MA. THEN MDW	WNN	370
C	MUST SATISFY MDW.GE.M. THE CONDITION MDW.LT.M	WNN	380
C	IS AN ERROR.	WNN	390
C		WNN	400
C	THE ARRAY W(*,*) CONTAINS THE MATRICES AND VECTORS	WNN	410
C		WNN	420
C	(E F)	WNN	430
C	(A B)	WNN	440
C		WNN	450
C	IN ROWS AND COLUMNS 1,...,M AND 1,...,N+1	WNN	460
C	RESPECTIVELY. COLUMNS 1,...,L CORRESPOND TO	WNN	470
C	UNCONSTRAINED VARIABLES X(1),...,X(L). THE	WNN	480
C	REMAINING VARIABLES ARE CONSTRAINED TO BE	WNN	490
C	NONNEGATIVE. THE CONDITION L.LT.0 .OR. L.GT.N IS	WNN	500
C	AN ERROR.	WNN	510
C		WNN	520
C	PRGOPT(*) THIS ARRAY IS THE OPTION VECTOR.	WNN	530
C	IF THE USER IS SATISFIED WITH THE NOMINAL	WNN	540
C	SUBPROGRAM FEATURES SET	WNN	550
C		WNN	560
C	PRGOPT(1)=1 (OR PRGOPT(1)=1.0)	WNN	570

ALGORITHM 588

Fast Hankel Transforms Using Related and Lagged Convolutions

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U.S. Geological Survey

Categories and Subject Descriptors: F.2.1 [Analysis of Algorithms and Problem Complexity]: Numerical Algorithms and Problems—*computation of transforms*; G.1.2 [Numerical Analysis]: Approximation—*linear approximation*; G.m [Mathematics of Computing]: Miscellaneous—*FORTRAN*

General Terms: Algorithms

Additional Key Words and Phrases: Hankel transforms of integer order, Bessel functions of the first kind, convolution integrals, linear digital filters

DESCRIPTION

Subprogram **HANKEL()**, for evaluating complex Hankel transforms, is a complement of [1], where the algorithm is discussed in more detail. A double-precision real version, subprogram **DHANKL()**, is also provided. The calling sequence for **DHANKL()** is ordered the same way as is described below for **HANKEL()**, except all **REAL** and **COMPLEX** declarations in **HANKEL()** are declared **DOUBLE PRECISION** in **DHANKL()**.

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ALGORITHM

[Summary information and a part of the listing is printed here. The entire package, including subprogram **HANKEL()** and three drivers, followed by **DHANKL()** and one driver, is available from the ACM Algorithms Distribution Service.]

Module Names

SUBROUTINE HANKEL()
[MAIN PROGRAM—HANKEL TEST 1]
COMPLEX FUNCTION FUNT1()
FUNCTION EXPM()
[MAIN PROGRAM—HANKEL TEST 2]
COMPLEX FUNCTION FUNT2()
FUNCTION SINHM()
[MAIN PROGRAM—HANKEL TEST 3]
COMPLEX FUNCTION FUNT3()

SUBROUTINE DHANKL()
[MAIN PROGRAM—DHANKL TEST 1D]

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ACM Transactions on Mathematical Software, Vol. 8, No. 4, December 1982, Pages 369-370 .

DOUBLE PRECISION FUNCTION DFUNT1()
 DOUBLE PRECISION FUNCTION DEXPM()

```

      SUBROUTINE HANKEL(BMAX,NB,NREL,TOL,NTOL,NORD,FUN1,IJREL,ZWORK,
      * ZANS,ARG,NOFUN1,IERR)
C=====
      INTEGER NB,NREL,NTOL,NORD(NREL),IJREL(2,NREL),NOFUN1,IERR
      REAL BMAX,TOL,ARG(NB)
      COMPLEX ZWORK(283,NREL),ZANS(NB,NREL)
C=====
C
C
C PURPOSE
C
C THE PURPOSE OF SUBPROGRAM HANKEL IS TO PROVIDE IN SINGLE PRECISION
C A GENERAL ALGORITHM FOR FAST COMPLEX HANKEL TRANSFORMS OF ORDERS
C 0 AND 1 USING RELATED AND LAGGED CONVOLUTIONS.
C
C AUTHOR
C
C ANDERSON, W.L., U.S. GEOLOGICAL SURVEY, DENVER, COLORADO.
C
C REFERENCES
C
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C TRANSFORMS OF ORDERS 0 AND 1 BY ADAPTIVE DIGITAL FILTERING.
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C
C LANGUAGE
C
C ANS-FORTRAN (X3.9-1966) IS USED, WITH THE EXCEPTION OF THE
C CHARACTERS F,!,&,.,1,J APPEARING IN SOME COMMENT STATEMENTS.

```

ALGORITHM 589

SICEDR: A FORTRAN Subroutine for Improving the Accuracy of Computed Matrix Eigenvalues

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1. INTRODUCTION

SICEDR is a FORTRAN subroutine for improving the accuracy of a computed real eigenvalue and improving or computing the associated eigenvector. It is first used to generate information during the determination of the eigenvalues by the Schur decomposition technique. In particular, the Schur decomposition technique results in an orthogonal matrix Q and an upper quasi-triangular matrix T , such that

$$A = QTQ^T. \quad (1.1)$$

Matrices A , Q , and T and the approximate eigenvalue, say λ , are then used in the improvement phase. **SICEDR** uses an iterative method similar to iterative improvement for linear systems to improve the accuracy of λ and improve or compute the eigenvector x in $O(n^2)$ work, where n is the order of the matrix A . The method used in **SICEDR** is described in [1, 5].

2. USAGE

For a description of the calling sequence, see the listing presented at the end of this paper.

SICEDR factors the matrix into its Schur decomposition, and this is termed the pre-**SICE** phase. A modification of the EISPACK routine **HQR2** is used for this purpose (see [4, pp. 100–101] for details).

During the improvement phase (or **SICE** phase), **SICEDR** is called with matrices A , T , and Q and the approximate eigenvalue W . On return from **SICEDR** the improved eigenvalue is stored in W and the improved or computed eigenvector in X . W and X contain the corrected eigenpair produced by the method at the next to last iteration. A still more accurate eigenpair can be formed if on return from **SICEDR** the user adds W to CW and X to CX in double

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precision and saves the results in double precision, where CW and CX contain the corrections to W and X , respectively, at the final iteration when the method terminates.

The routine **SICEDR** is provided as a driver to simplify usage. **SICEDR** performs both the pre-**SICE** phase (reduction to upper form) and the **SICE** phase (improvement).

In addition to the EISPACK package, the LINPACK [2] routine **STRSL** and the BLAS [3] package are used to perform fundamental operations.

To produce a double-precision version of **SICEDR** in addition to making the obvious changes, such as changing variables to double precision and using double-precision intrinsic functions, it is necessary to replace routine **STRSL** by **DTRSL** and use the double-precision version of the BLAS. In addition, there are two critical parts of the calculation which *must* be performed in an extended precision. In the single-precision version, the routine **SDDDOT** and **SDADD** perform double-precision accumulations. In a double-precision version, they must be replaced by, say **DQDDOT** and **DQADD**, the extended precision counterparts. These extended precision routines would accumulate in quadruple precision, a nonstandard FORTRAN construct. Such routines are not included here, but can be constructed from the BLAS routines **DQDOTI** and **DQDOTA**.

3. SUMMARY OF THE ALGORITHM

We begin with a brief discussion of the basic algorithm as described in [1].

If λ , x is an approximate eigenpair, and $\lambda + \mu$, $x + \tilde{y}$ is a neighboring eigenpair, then

$$A(x + \tilde{y}) = (\lambda + \mu)(x + \tilde{y}), \quad (3.1)$$

this relation being exact. We assume that x is normalized so that $\|x\|_\infty = 1 = x_s$, and we remove the degree of arbitrariness in \tilde{y} by requiring that $\tilde{y}_s = 0$. Rearranging this equation, we have

$$(A - \lambda I)\tilde{y} - \mu x = \lambda x - Ax + \mu\tilde{y}, \quad (3.2)$$

where the last term on the right will be of second order in the errors of λ , x . The equation above may be simplified by the introduction of a vector y defined by

$$y^T = (y_1, y_2, \dots, y_{s-1}, \mu, y_{s+1}, \dots, y_n), \quad (3.3)$$

so that y gives the full information on both μ and \tilde{y} . The above equation then becomes

$$By = r + y_s \tilde{y}, \quad (3.4)$$

where $r = \lambda x - Ax$ is the residual vector corresponding to λ , x and B is the matrix $A - \lambda I$ with column s replaced by $-x$.

When the original approximate eigenpairs have been found by Francis' double **QR** algorithm, we have an orthogonal matrix Q and a quasi upper triangular matrix T , such that

$$A = QTQ^T. \quad (3.5)$$

T is triangular apart from possible 2×2 diagonal blocks corresponding to complex conjugate eigenvalues. We solved a succession of linear systems of the form (3.4) with varying right-hand sides g . It will be convenient to introduce the generic notation $Z - \lambda I = Z_\lambda$ and $(Z - \lambda I)e_s = Z_\lambda e_s = z_{\lambda s}$. We may then rewrite eq. (3.4) in the form

$$[A_\lambda - (x + a_{\lambda s})e_s^T]y = (A_\lambda + ce_s^T)y = g, \quad (3.6)$$

where

$$c = -x - a_{\lambda s}.$$

Using the orthogonal factorization, we have

$$Q[T_\lambda + Q^T c e_s^T Q] Q^T y = g,$$

giving

$$(T_\lambda + d f^T) Q^T y = Q^T g, \quad (3.7)$$

where

$$d = Q^T c, \quad f^T = e_s^T Q \quad \text{and} \quad g = r + y_s \tilde{y}.$$

The matrix $d f^T$ is a rank one modification of the quasi-triangular matrix T_λ . To solve this system, we need to retriangularize $T_\lambda + d f^T$. Accordingly, we premultiply the system by two orthogonal matrices Q_1 and Q_2 , giving

$$Q_2 Q_1 (T_\lambda + d f^T) Q^T y = Q_2 Q_1 Q^T g, \quad (3.8)$$

where Q_1 and Q_2 are products of elementary plane rotations determined as follows.

The matrix Q_1 is such that

$$Q_1 d = (P_2 P_3 \cdots P_n) d = \gamma e_1 \quad \text{where} \quad \gamma = \|d\|_2 \quad (3.9)$$

and P_i is a rotation in the $(i-1, i)$ plane designed to annihilate the i th component of $P_{i+1} P_{i+2} \cdots P_n d$. We have

$$Q_1 (T_\lambda + d f^T) = Q_1 T_\lambda + \gamma e_1 f^T, \quad (3.10)$$

where $Q_1 T_\lambda$ is upper Hessenberg, while $\gamma e_1 f^T$ is null except in the first row. Hence the right-hand side is also an upper Hessenberg matrix H . H may now be reduced to upper triangular form, \bar{T}_λ , by premultiplication with Q_2 defined by

$$Q_2 = P'_n \cdots P'_3 P'_2, \quad (3.11)$$

where the premultiplication by P'_i annihilates the element $(i, i-1)$ of the current matrix by a rotation in the $(i-1, i)$ plane. Hence, the triangular system remains to be solved

$$\bar{T}_\lambda Q^T y = Q_2 Q_1 Q^T g. \quad (3.12)$$

A system with the matrix B may thus be solved in $O(n^2)$ operations.

By its nature, eq. (3.2), which leads to eq. (3.12), is mildly nonlinear. Thus we repeat the process with $x + \tilde{y}$ and $\lambda + \mu$ as the new approximations. The convergence theorem for this iterative procedure can be found in [1]. Since an orthogonal triangularization of A is available, it becomes practical to update the matrix B at each stage of the iteration, using the current approximation to the eigenpair. Accordingly, we treat the $(p+1)$ th step of the iteration as though it were the first step in the basic iteration starting with values $\lambda^{(p)}$ and $x^{(p)}$. Since we treat each iterate as if it were the first, the term $\mu \tilde{y}$ in eq. (3.2) is zero. Thus $g = \lambda x - Ax$. The algorithm then becomes

$$(A - \lambda^{(p)} I) \tilde{y}^{(p)} - y_s^{(p)} \tilde{x}^{(p)} = r^{(p)} = (\lambda^{(p)} I - A) x^{(p)} \quad (3.13)$$

where

$$x^{(p+1)} = x^{(p)} + \tilde{y}^{(p)}, \quad \lambda^{(p+1)} = \lambda^{(p)} + y_s^{(p)}.$$

We may rewrite this as

$$B^{(p)} y^{(p)} = r^{(p)}, \quad (3.14)$$

where

$$B^{(p)} = A - \lambda^{(p)} I + c^{(p)} e_s^T, \quad c^{(p)} = -x - \alpha_{\lambda_s}^{(p)}.$$

We now write

$$B^{(p)} = Q(T_\lambda^{(p)} + Q^T c^{(p)} e_s^T Q) Q^T$$

$$= Q(T_{\lambda}^{(p)} + d^{(p)}f^T)Q^T \quad (3.15)$$

and solve (3.14) by the triangular system

$$\bar{T}_{\lambda}^{(p)}Q^T y^{(p)} = Q_{\frac{1}{2}}^{(p)}Q_1^{(p)}Q^T r^{(p)}. \quad (3.16)$$

Convergence is detected by examining successive values for the correction to the eigenvalue. When the previous correction is smaller by a factor of 2 than the current correction, the iteration is stopped.

Note that Q and f will be independent of p if s is not changing from one iteration to the next. The rotations involved in $Q_1^{(p)}$ and $Q_{\frac{1}{2}}^{(p)}$, on the other hand, differ from one iteration to the next; but because the number of operations in each retriangularization is $O(n^2)$ and since some $\frac{1}{2}n^2$ multiplication and additions are necessarily involved in the solution of a triangular system, this is quite acceptable.

If we examine the number of operations for the reduction to quasi-triangular form, the *pre-SICE* phase, we see that there are roughly $10n^3 + 30n^2$ operations involved (here, an operation means a multiplication followed by an addition). For the *SICE* phase there are approximately $13n^2$ operations per iteration. Assuming 3 iterations to improve an eigenpair, the total count is approximately $39n^2$ operations.

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ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```

SUBROUTINE SICEDR(LD, N, A, T, Q, W, X, CW, CX, K, Y1, Y2, R, Z1, SIC 10
* WR, WI, JOB, INFO, INIT) SIC 20
C SIC 30
INTEGER I, J, IPI, IERR, LD, N, K, JOB, INFO, INIT SIC 40
REAL A(LD,1), T(LD,1), Q(LD,1), WR(1), WI(1) SIC 50
REAL W, X(1), CW, CX(1), Y1(1), Y2(1), R(1), Z1(1) SIC 60
C SIC 70
C THIS SUBROUTINE PROVIDES TWO FUNCTIONS DEPENDING ON SIC 80
C THE VALUE OF INIT. SIC 90
C SIC 100
C 1) TO REDUCE A MATRIX TO QUASI-TRIANGULAR FORM, SIC 110
C ACCUMULATING THE ORTHOGONAL TRANSFORMATIONS, SIC 120
C AND DETERMINE THE EIGENVALUES OF THE MATRIX. SIC 130
C SIC 140
C 2) TO IMPROVE THE ACCURACY OF AN EIGENVALUE AND SIC 150
C IMPROVE THE EIGENVECTOR OR COMPUTE AND IMPROVE THE SIC 160
C EIGENVECTOR, GIVEN THE ORIGINAL MATRIX, AN SIC 170
C APPROXIMATE EIGENVALUE, THE QUASI-TRIANGULAR MATRIX, SIC 180
C AND THE ORTHOGONAL MATRIX WHICH PRODUCED THE SIC 190
C QUASI-TRIANGULAR FORM. SIC 200
C SIC 210
C THESE FUNCTIONS ARE SIGNALLED BY THE PARAMETER INIT. SIC 220

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ALGORITHM 590

DSUBSP and EXCHQZ: FORTRAN

Subroutines for Computing Deflating Subspaces with Specified Spectrum

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1. DESCRIPTION

A reliable and widely available method to compute the generalized eigenvalues of an $n \times n$ real regular (i.e., invertible) pencil $\lambda B - A$ is the so-called *QZ*-algorithm [1, 3]. This algorithm constructs orthogonal row and column transformations Q_1 and Z_1 such that the transformed pencil

$$\lambda B_1 - A_1 = Q_1(\lambda B - A)Z_1 \quad (1)$$

is in “quasi-triangular” form, that is, with B_1 in upper triangular form and A_1 in block triangular Hessenberg form with 1×1 and 2×2 diagonal blocks, as illustrated below:

$$B_1 = \begin{bmatrix} x & x & x & x & x & x \\ & x & x & x & x & x \\ & & x & x & x & x \\ & & & x & x & x \\ & & & & x & x \\ & & & & & x \end{bmatrix}, \quad A_1 = \begin{bmatrix} x & x & x & x & x & x \\ & x & x & x & x & x \\ & & x & x & x & x \\ & & & x & x & x \\ & & & & x & x \\ & & & & & x & x \end{bmatrix}. \quad (2)$$

The 1×1 diagonal pencils of $\lambda B_1 - A_1$ contain the real generalized eigenvalues of the pencil $\lambda B - A$, and the 2×2 diagonal pencils of $\lambda B_1 - A_1$ contain the complex generalized eigenvalues, a conjugate pair to each 2×2 pencil.

The FORTRAN subroutines **DSUBSP** and **EXCHQZ** allow one to update the decomposition (1) by postmultiplying Z_1 by Z_2 and premultiplying Q_1 by Q_2 such that

$$\lambda B_2 - A_2 = Q_2 Q_1 (\lambda B - A) Z_1 Z_2 \quad (3)$$

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is still in quasi-triangular form, but in addition, a specific ordering of generalized eigenvalues is obtained in this form. This is important in several applications [6, 7]. Indeed, when the generalized eigenvalues $\lambda_1, \dots, \lambda_n$ are ordered such that $\lambda_1, \dots, \lambda_k$ are inside a region Γ and $\lambda_{k+1}, \dots, \lambda_n$ are outside this region, then the space spanned by the first k orthonormal columns of $Z = Z_1 Z_2$ is the deflating subspace [4] of $\lambda B - A$ corresponding to the spectrum inside Γ . In several applications, deflating subspaces have to be computed for different regions Γ . Of course, Γ has to be symmetric with respect to the real axis, since complex pairs of eigenvalues have to be categorized both outside or inside Γ in order for this deflating subspace to be real. No further assumptions have to be imposed on Γ : it can be open or not, connected or not.

The user has to provide a function describing the region Γ by testing whether the generalized eigenvalues of a 1×1 or 2×2 (diagonal) pencil lies inside or outside the region Γ . This function must be of the type

INTEGER FUNCTION FTEST (LSIZE, ALPHA, BETA, S, P)

with parameters:

- LSIZE** An integer containing the size of the considered pencil (1 or 2).
- ALPHA,BETA** Two real variables. In case **LSIZE**=1, the generalized eigenvalue of the considered pencil is given by **ALPHA/BETA**, which may be infinite when **BETA**=0.
- S, P** Two real variables. In case **LSIZE**=2, they contain the sum and product of the two complex conjugate generalized eigenvalues of the considered pencil.
- FTEST** The function value, which is put equal to 1 when the generalized eigenvalue(s) of the considered pencil is (are) inside the specified region Γ , and equal to -1 otherwise.

Simple examples for such routines are given by the functions **FIN**, **FOUT**, **FOLHP**, and **FCRHP**, describing the regions inside and outside the unit circle, the open left half-plane, and the closed right half-plane, respectively. Their listings are included below as templates for the user.

This routine is then used as parameter for the subroutine **DSUBSP**, which reorders the 1×1 and 2×2 diagonal pencils of the quasi-triangular form $\lambda B_1 - A_1$ such that those with generalized eigenvalues inside Γ appear first. The calling sequence for **DSUBSP** is

CALL DSUBSP (NMAX, N, A, B, Z, FTEST, EPS, NDIM, FAIL, IND)

with (parameters preceded by an asterisk are altered by the subroutine):

- NMAX** An integer containing the first dimension of the arrays **A**, **B**, and **Z**.
- N** An integer containing the current order of **A**, **B**, and **Z**.
- ;axA,*B** Doubly subscripted real arrays containing the pencil to be reordered. On return, $\lambda B - A$ contains the final quasi-triangular pencil, rearranged with respect to the region Γ specified by **FTEST**.
- *Z** A doubly subscripted real array into which the reducing column transformation is postmultiplied.
- FTEST** The integer function provided by the user to describe the region Γ of interest.
- EPS** A real number used as the convergence criterion. Maximal accuracy is obtained when **EPS** is set equal to $\text{relpr} \times \max\{\|\mathbf{A}\|_2, \|\mathbf{B}\|_2\}$, where **relpr** is the relative precision of the computer used. Smaller values of **EPS** will increase the amount of work without significantly improving the accuracy.
- *NDIM** An integer giving the dimension of the computed deflating subspace.
- *FAIL** A logical variable that on normal return is **.FALSE**. If the iterative part of the algorithm does not converge, **FAIL** is set to **.TRUE**.

*IND An integer working array of dimension at least N.

DSUBSP is to be used together with the EISPACK programs **QZHES**, **QZIT**, and **QZVAL** [1] to reduce a full pencil $\lambda B - A$ to quasi-triangular form with the eigenvalues inside the contour Γ appearing first on diagonal. (For the explanation of the parameters **EPS1**, **IERR**, **ALPHAR**, **ALPHAI**, and **BETA** in these routines, see [1].):

CALL **QZHES** (NMAX, N, A, B, .TRUE., Z)
 CALL **QZIT** (NMAX, N, A, B, EPS1, .TRUE., Z, IERR)
 CALL **QZVAL** (NMAX, N, A, B, ALPHAR, ALPHAI, BETA, .TRUE., Z)
 CALL **DSUBSP** (NMAX, N, A, B, Z, FTEST, EPS, NDIM, FAIL, IND)

Besides the function **FTEST**, describing the region Γ of interest, **DSUBSP** also uses the subroutines **EXCHQZ**, **GIV**, and **SROT**. **EXCHQZ** is a FORTRAN subroutine to interchange two adjacent (1×1 and/or 2×2) pencils of a quasi-triangular form. Specifically, it is supposed that A has a block of order l_1 starting at the l th diagonal element, and a block of order l_2 starting at the $(l + l_1)$ th diagonal element (illustrated below for $n = 5$, $l = 2$, $l_1 = 2$, $l_2 = 1$):

$$B = \begin{bmatrix} x & x & x & x & x \\ & x & x & x & x \\ & & x & x & x \\ & & & x & x \\ & & & & x \end{bmatrix}, \quad A = \begin{bmatrix} x & x & x & x & x \\ & x & x & x & x \\ & & x & x & x \\ & & & x & x \\ & & & & x \end{bmatrix}. \quad (4)$$

EXCHQZ constructs orthogonal row and column transformations V and W such that $V(\lambda B - A)W$ has consecutive blocks of order l_2 and l_1 at the l th diagonal element (illustrated for the example above):

$$VBW = \begin{bmatrix} x & x & x & x & x \\ & x & x & x & x \\ & & x & x & x \\ & & & x & x \\ & & & & x \end{bmatrix}, \quad VAW = \begin{bmatrix} x & x & x & x & x \\ & x & x & x & x \\ & & x & x & x \\ & & & x & x \\ & & & & x \end{bmatrix}. \quad (5)$$

The generalized eigenvalues associated with each diagonal block are interchanged along with the blocks. The column transformation W is postmultiplied into the array Z ; the row transformation V is not stored since it is not required in the computation of bases of deflating subspaces.

The calling sequence for **EXCHQZ** is

CALL **EXCHQZ** (NMAX, N, A, B, Z, L, LS1, LS2, EPS, FAIL)

with (parameters preceded by an asterisk are altered by the subroutine):

NMAX An integer containing the first dimension of **A**, **B**, and **Z**.

N An integer containing the current order of **A**, **B**, and **Z**.

***A**, ***B** Doubly subscripted real arrays containing the pencil to be reordered.

***Z** A doubly subscripted real array into which the updating column transformation is postmultiplied.

L An integer containing the leading diagonal position of the first block to be interchanged.

LS1 An integer containing the size of the first block.

- LS2** An integer containing the size of the second block.
- EPS** A convergence criterion (cf. **EPS** in the calling sequence of **DSUBSP**).
- *FAIL** A logical variable that on normal return is **.FALSE.**. If the iterative part of the algorithm does not converge, **FAIL** is set to **.TRUE.**.

EXCHQZ requires the subroutines **GIV** and **SROT**, which are elementary routines that construct and perform Givens rotations on column or rows. **SROT** is a BLA subroutine and **GIV** is a modification of **SROTG** of the BLAS package [2] in order to allow a more compact code for **EXCHQZ**, which contains several calls to these routines.

2. METHOD AND PROGRAMMING DETAILS

The subroutine **DSUBSP** makes a first pass through the diagonal blocks of the quasi-triangular form $\lambda B - A$ in order to determine the sizes of the diagonal blocks and the locations of their generalized eigenvalues with respect to Γ . During this pass the integer vector **IND**(•) is created with entries ± 1 or ± 2 . Here **sign** [**IND**(**I**)] refers to the location of the generalized eigenvalues of block **I** with respect to Γ (+ for inside Γ , - for outside Γ) and **abs**[**IND**(**I**)] indicates the size of this block. The entries of this vector are then rearranged such that the plus signs appear first. This is done using a "bubble sort," that is, each time a plus sign follows a minus sign, it is moved in front of all its preceding minus signs via consecutive permutations. Each permutation, of course, involves an interchanging of two consecutive blocks using the subroutine **EXCHQZ**.

EXCHQZ works in a fashion similar to the routine **EXCHNG**, developed for the reordering of the standard eigenvalue problem [5] (this is also the reason why an apparent parallelism with that paper has been pursued here). To interchange two consecutive blocks where at least one of them has order 1, a shift is performed to the real eigenvalue of a 1×1 block (in [5], this was only done for the interchange of two 1×1 blocks). Givens rotations are then constructed in a straightforward manner to interchange the shifted zero eigenvalue to the other block. The construction of the Givens transformations needed for the exchange of the blocks is done in such a way as to ensure backward stability.

In the case of two 2×2 blocks, an arbitrary **QZ**-step is performed on both blocks in order to eliminate the uncoupling between them. Then a sequence of **QZ**-steps using a previously determined shift is performed on both blocks. A decoupling with the blocks in the desired order is usually obtained with a few steps (rarely more than two). If within 30 iterations no decoupling is obtained, the subroutine gives an error return. The criterion used is that the coupling element of the two blocks in the Hessenberg form of A is smaller than **EPS**. Since it does not make sense to force this coupling element to be smaller than the errors in the rest of the pencil, one should not choose **EPS** smaller than $\text{relpr} \times \max\{\|A\|_2, \|B\|_2\}$, where **relpr** is the relative precision of the computer used. A good choice for **EPS** is the "estimated" absolute precision of the (sometimes measured) data in A and B . Since the pencils used here are of dimension 4×4 , the **QZ** steps are implemented with Givens transformations instead of Householder transformations, which turns out to be economical. More details are given in [7], where a proof of the backward stability of the method is also given.

EXCHQZ uses the routines **GIV** and **SROT**, which construct a 2×2 Givens rotation to zero out an element of a 2-vector, and perform it, respectively, on two columns or rows of a specified matrix.

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ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Algorithms Distribution Service.]

```

SUBROUTINE DSUBSP(NMAX, N, A, B, Z, FTEST, EPS, NDIM, FAIL, IND) DSU 10
INTEGER NMAX, N, FTEST, NDIM, IND(N) DSU 20
LOGICAL FAIL DSU 30
REAL A(NMAX,N), B(NMAX,N), Z(NMAX,N), EPS DSU 40
C* DSU 50
C* GIVEN THE UPPER TRIANGULAR MATRIX B AND UPPER HESSENBERG MATRIX A DSU 60
C* WITH 1X1 OR 2X2 DIAGONAL BLOCKS, THIS ROUTINE REORDERS THE DIAGONAL DSU 70
C* BLOCKS ALONG WITH THEIR GENERALIZED EIGENVALUES BY CONSTRUCTING EQUI-DSU 80
C* VALENCE TRANSFORMATIONS QT AND ZT. THE ROW TRANSFORMATION ZT IS ALSO DSU 90
C* PERFORMED ON THE GIVEN (INITIAL) TRANSFORMATION Z (RESULTING FROM A DSU 100
C* POSSIBLE PREVIOUS STEP OR INITIALIZED WITH THE IDENTITY MATRIX). DSU 110
C* AFTER REORDERING, THE EIGENVALUES INSIDE THE REGION SPECIFIED BY THE DSU 120
C* FUNCTION FTEST APPEAR AT THE TOP. IF NDIM IS THEIR NUMBER THEN THE DSU 130
C* NDIM FIRST COLUMNS OF Z SPAN THE REQUESTED SUBSPACE. DSUBSP REQUIRES DSU 140
C* THE SUBROUTINE EXCHQZ AND THE INTEGER FUNCTION FTEST WHICH HAS TO BE DSU 150
C* PROVIDED BY THE USER. THE PARAMETERS IN THE CALLING SEQUENCE ARE : DSU 160
C* (STARRED PARAMETERS ARE ALTERED BY THE SUBROUTINE) DSU 170
C* DSU 180
C* NMAX THE FIRST DIMENSION OF A, B AND Z DSU 190
C* N THE ORDER OF A, B AND Z DSU 200
C* *A,*B THE MATRIX PAIR WHOSE BLOCKS ARE TO BE REORDERED. DSU 210
C* *Z UPON RETURN THIS ARRAY IS MULTIPLIED BY THE COLUMN DSU 220
C* TRANSFORMATION ZT. DSU 230
C* FTEST(LS,ALPHA,BETA,S,P) AN INTEGER FUNCTION DESCRIBING THE DSU 240
C* SPECTRUM OF THE DEFLATING SUBSPACE TO BE COMPUTED: DSU 250
C* WHEN LS=1 FTEST CHECKS IF ALPHA/BETA IS IN THAT SPECTRUM DSU 260
C* WHEN LS=2 FTEST CHECKS IF THE TWO COMPLEX CONJUGATE DSU 270
C* ROOTS WITH SUM S AND PRODUCT P ARE IN THAT SPECTRUM DSU 280
C* IF THE ANSWER IS POSITIVE, FTEST=1, OTHERWISE FTEST=-1 DSU 290
C* EPS THE REQUIRED ABSOLUTE ACCURACY OF THE RESULT DSU 300
C* *NDIM AN INTEGER GIVING THE DIMENSION OF THE COMPUTED DSU 310
C* DEFLATING SUBSPACE DSU 320
C* *FAIL A LOGICAL VARIABLE WHICH IS FALSE ON A NORMAL RETURN, DSU 330
C* TRUE OTHERWISE (WHEN EXCHQZ FAILS) DSU 340
C* *IND AN INTEGER WORKING ARRAY OF DIMENSION AT LEAST N DSU 350
C* DSU 360
INTEGER L, LS, LS1, LS2, L1, LL, NUM, IS, L2I, L2K, I, K, II, DSU 370
* ISTEP, IFIRST DSU 380
REAL S, P, D, ALPHA, BETA DSU 390
FAIL = .TRUE. DSU 400
NDIM = 0 DSU 410
NUM = 0 DSU 420
L = 0 DSU 430
LS = 1 DSU 440
C*** CONSTRUCT ARRAY IND(I) WHERE : DSU 450
C*** IABS(IND(I)) IS THE SIZE OF THE BLOCK I DSU 460
C*** SIGN(IND(I)) INDICATES THE LOCATION OF ITS EIGENVALUES DSU 470
C*** (AS DETERMINED BY FTEST). DSU 480
C*** NUM IS THE NUMBER OF ELEMENTS IN THIS ARRAY DSU 490
DO 30 LL=1,N DSU 500
L = L + LS DSU 510

```

COLLECTED ALGORITHMS (cont.)

590-P 6- 0

IF (L.GT.N) GO TO 40	DSU 520
L1 = L + 1	DSU 530
IF (L1.GT.N) GO TO 10	DSU 540
IF (A(L1,L).EQ.0.) GO TO 10	DSU 550
C* HERE A 2X2 BLOCK IS CHECKED *	DSU 560
LS = 2	DSU 570
D = B(L,L)*B(L1,L1)	DSU 580
S = (A(L,L)*B(L1,L1)+A(L1,L1)*B(L,L)-A(L1,L)*B(L,L1))/D	DSU 590
P = (A(L,L)*A(L1,L1)-A(L,L1)*A(L1,L))/D	DSU 600
IS = FTEST(LS,ALPHA,BETA,S,P)	DSU 610
GO TO 20	DSU 620
C* HERE A 1X1 BLOCK IS CHECKED *	DSU 630
10 LS = 1	DSU 640
IS = FTEST(LS,A(L,L),B(L,L),S,P)	DSU 650
20 NUM = NUM + 1	DSU 660
IF (IS.EQ.1) NDIM = NDIM + LS	DSU 670
IND(NUM) = LS*IS	DSU 680
30 CONTINUE	DSU 690
C*** REORDER BLOCKS SUCH THAT THOSE WITH POSITIVE VALUE	DSU 700
C*** OF IND(.) APPEAR FIRST.	DSU 710

ALGORITHM 591

A Comprehensive, Matrix-Free Algorithm for Analysis of Variance

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Categories and Subject Descriptors: G.3 [Mathematics of Computing]: Probability and Statistics—*statistical computing; statistical software*; G.m [Mathematics of Computing]: Miscellaneous—*FORTRAN*

General Terms: Algorithms

Additional Key Words and Phrases: Linear models, analysis of variance, unbalanced data, missing cells, estimation, hypothesis testing, storage-efficient algorithm, iterative improvement, balanced data operators, algebraic-model specification, rank computations, G-inverse solution.

1. INTRODUCTION

Some of the work to be described is reminiscent of **AARDVARK** [5], an analysis-of-variance package, at Iowa State University in the early 1960s. This program was used for most of the large number of analyses processed at Iowa State at that time; however, the program was definitely not transportable, being written in a mix of assembler and FORTRAN for a nonstandard IBM 7074. Subsequent versions, lacking some of its initial capabilities but adding others, were prepared for other machines, including the IBM 360 (see for example [7]). One distinctive feature of the program that greatly aided in its usability was algebraic specification of the statistical model by the user [14], [18]. This facility has now become common in statistical packages.

The initial **AARDVARK** relied principally on balanced analysis-of-variance (AOV) algorithms and approximate statistical methods were applied to treat unbalanced data. An iterative AOV algorithm was later developed by Hemmerle [12], which permits using balanced analysis-of-variance algorithms to obtain exact statistical results for unbalanced data. This algorithm along with subsequent work will be exploited in our development of a comprehensive analysis-of-variance algorithm for balanced or unbalanced data.

Balanced analysis-of-variance algorithms frequently represent a level of elegance of logic that is difficult to surpass. One fine example is the algorithm due to H. O. Hartley [3], which obtains a complete factorial decomposition. Another is Wilkinson's recursive algorithm [23]. Unfortunately, these algorithms are also frequently difficult to explain using straightforward notation; they are likely to involve matrices with complicated structures implicitly rather than explicitly.

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Balanced analysis-of-variance algorithms normally do not require the coding of indicator variables, either externally or internally by the program. They possess the virtue of efficiency, both with respect to the number of operations performed and the amount of storage required.

Data, on the other hand, are frequently unbalanced and are likely to have missing cells. Both theoretically and computationally, the order of complexity of analysis-of-variance problems increases down the following categories:

- (1) balanced data;
- (2) unbalanced data with no missing cells;
- (3) unbalanced data with missing cells.

In some cases, the unbalanced data arise from a balanced design in which some of the observations are missing as a result of the experiment. Computationally, there are a number of missing data algorithms, notably Healy and Westmacott's iterative algorithm [4] and Rubin's noniterative algorithm [16], which essentially fill in the missing values for use with balanced algorithms. One could also use dummy covariates for each missing cell and then compute a balanced analysis of covariance. These approaches have the disadvantage that, in general, they require balancing all of the data (including the missing observations) to have equal cell frequencies.

Almost all algorithms which seek to process unbalanced data with any degree of generality in an analysis-of-variance situation, including those in statistical packages, are matrix-oriented least squares or regression algorithms. These involve the creation of indicator variables, to cast the analysis-of-variance model as a regression model, and the formation of the design matrix X_0 or the coefficients matrix of the normal equations $X_0'X_0$. The normal equations are then solved using either orthogonalization or elimination methods. In some cases, SAS Procedure GLM [1], for example, an explicit G-inverse is obtained for $X_0'X_0$. One problem with the matrix-oriented approach is that the dimensions of X_0 or $X_0'X_0$ can be overwhelming for moderate to large analysis-of-variance problems. This approach is also extremely inefficient when the data are balanced.

2. FEATURES AND DESIGN OBJECTIVES

In what follows, we describe a global algorithm for the analysis of variance with the following features.

(1) Balanced data, unbalanced data, and unbalanced data with missing cells are all processed by the algorithm. This is accomplished without losing the operational efficiencies obtainable from balanced data and without applying approximate statistical methods to the unbalanced data.

(2) The algorithm is very general with respect to the kind of problems it can handle. Specifically, it bases its calculations upon an algebraically specified analysis-of-variance model of the type discussed in Searle [19]. This includes models with crossed factors, nested factors, and interactions between factors. This general model, along with the facility to handle missing cells, also includes such designs as incomplete blocks, lattices, and Latin squares.

(3) Very large problems may be processed using a relatively small amount of computer storage. With one minor exception, no matrices are stored and no explicit matrix operations are performed. In particular, neither X_0 nor $X_0'X_0$ is stored or computed. An exact G-inverse solution to the normal equations is obtained without ever computing a G-inverse. The rank of the design matrix X_0 is obtained from the pattern of missing cells without explicit operations on X_0 .

(4) The algorithm provides a heuristic optimum of maximizing analysis-of-variance capabilities while minimizing lines of source code. Since a principal attribute of the algorithm is its limited demands for array storage, an effort was made to also constrain program storage through a judicious selection of functional capabilities and the production of tight code. Consequently, the algorithm should be well suited for use on small computers.

(5) The user is given reasonable facilities in specifying his analysis-of-variance problem algebraically and various options alphanumerically in an essentially format-free manner.

(6) The algorithm is well suited to interactive usage. It was developed using **CALL/OS** at the University of Rhode Island and **TSO** (time-sharing option) at North Carolina State. It may, of course, also be used as a batch program. A mode parameter, interactive or batch, determines the nature of prompting, I/O, and error returns.

(7) The program has been written in American National Standard Basic FORTRAN for portability. The program has been verified by the Bell Laboratories PFORT verifier [17] and should run on any computer supported with a FORTRAN 66 compiler. With the exception of using the Hollerith data type in DATA statements, the program is also compatible with FORTRAN 77.

Along with the effort to constrain program storage, steps were taken to reduce the number of *lines* of source code. Lines of source (including **COMMENT** statements) can be a limiting factor with respect to file storage in interactive systems. Most of the following measures served to reduce program storage as well as lines of source code: limiting the number of specification statements, **WRITE** statements, **FORMAT** statements, and unnecessary **CONTINUE** statements; limiting the amount of constant alphanumeric information in **FORMAT** statements; using selective, concise **COMMENT** statements; accepting trivial operational inefficiencies which reduce lines of source code; avoiding modularity not contributing to operational efficiency or storage reduction; and using single-character identifiers when possible. The portability objective essentially limits one to storage of a single alphanumeric character per alphanumeric constant or array location. Rather than use additional code to interpret multiple-character identifiers, for such things as specified options, we limit these to one character.

As a consequence of these measures, the complete program for the global algorithm consists of 836 FORTRAN statements exclusive of **COMMENT** statements. The object code produced by the **WATFIV** compiler was contained in 36,656 bytes. In addition, the program was written in such a manner that several of the functional capabilities may easily be deleted, if necessary, to relax program storage.

3. CONSTITUENT ALGORITHMS

We have called the algorithm under discussion a global algorithm, since it encompasses the use of several independent algorithms. The principal independent algorithms, which comprise the global algorithm, are the following.

3.1 The Iterative AOV Algorithm

This algorithm iteratively applies expectation and residual operators (E/R operators) for balanced data to obtain exact results for unbalanced data. In other words, stated simply, the algorithm computes an analysis of variance for unbalanced data by successive computation of balanced analyses. Since the algorithm uses balanced E/R operators to obtain its results, there is no need to create indicator variables or form the design matrix X_0 or $X_0'X_0$. The algorithm operates upon a vector of cell sums and a vector of cell frequencies rather than upon a (potentially large) matrix so that array storage requirements are minimal. Missing cells are handled routinely and convergence of the algorithm is guaranteed [12]. A useful monotonic property serves to limit iteration when testing hypotheses. A balanced analysis of variance is a special case requiring only one iteration.

The basic algorithm [12] is very general and applies to any analysis of variance; however, one must know or determine the balanced E/R operators. It is in this context that one can apply much of the algorithmic lore associated with balanced analyses. An initial implementation using a relatively comprehensive class of balanced E/R operators was discussed by Hemmerle and Piette [6]. Many of the features of the global algorithm were included in this implementation; however,

certain unnecessary rank restrictions were placed upon the design matrix X_0 . The latter restrictions were subsequently removed by Hemmerle [10], who showed that with proper parameter selection the iterative AOV algorithm would obtain a G-inverse solution to the normal equations for any design matrix. An improved iterative approximation, which halves the number of iterations in testing hypotheses, is also given in [10] and is used here.

An iterative analysis of covariance algorithm was also developed by Hemmerle [11] as an extension of the iterative AOV algorithm. Although this covariance algorithm adapts well to the logic of the global algorithm and is array storage efficient, covariance capabilities were excluded here principally because of the additional program storage requirements.

3.2 An Iterative Rank Algorithm

With unbalanced data, including missing cells, the ranks of the design matrix and restricted design matrix must be computed to determine degrees of freedom for F statistics. Algorithms have been included (Section 3.5) to determine rank noniteratively from the pattern of missing cells for a given model when possible. When this cannot be accomplished, a recent adaptation of the iterative AOV algorithm is used to compute rank in one of two ways, depending upon the number of missing cells. Either way will produce monotonically increasing convergence to the integer rank value. For a small number of missing cells we use some matrix storage (the exception cited earlier) and obtain monotonically increasing quadratic convergence. Both methods are outlined in a subsequent section. A complete discussion is given in [9].

3.3 A Balanced Factorial Decomposition Algorithm

The technique for computing a balanced analysis of variance based upon a factorial decomposition of linear combinations of classifications means dates back to H. O. Hartley [3].

We use the algorithm described in Hemmerle [13, pp. 180-185] to obtain this decomposition. A one-to-one correspondence is maintained between the distinct arrays in this decomposition and the locations of an E/R list which we shall discuss shortly.

3.4 A Pooling Algorithm

In order to construct the E/R operators needed for the iterative AOV algorithms, we must "pool" (add together in the proper manner) distinct arrays in the factorial decomposition as dictated by the particular analysis-of-variance model. The complete factorial decomposition will be stored in a single linear, one-dimensional array, and we must be able to map one of the distinct arrays within this decomposition into another. We use essentially the same algorithm described in Schlater and Hemmerle [18] to do the mapping and pooling; however, we pool arrays of the decomposition rather than classification means, as do the latter.

3.5 Algorithms to Recognize Balance and Restructure Data

A balanced analysis of variance is a special case for the iterative AOV algorithm. When properly executed, final results will be obtained on the first iteration. Furthermore, iterative rank computations are unnecessary. With unbalanced data, there are also frequent situations in which some of the results may be obtained noniteratively. An example would be a full factorial model; the rank of the design matrix, the sum of squares for error (SSE), the sum of squares for regression (SSR), and a G-inverse solution to the normal equations may all be obtained from a balanced analysis of variance on cell means [12]. Also, when testing an hypothesis involving unbalanced data, the data structure for the model restricted by the hypothesis frequently has fewer dimensions than it has for the full model. A simple restructuring of the data may facilitate obtaining results without iterating. Essentially, this involves creating a surrogate vector of cell

frequencies based upon reduced dimensions for the same data; however, duplicate entries are made in this surrogate vector to maintain a proper one-to-one correspondence with the original vector of cell sums. Several algorithms have been incorporated into the global algorithm to recognize balance and restructure data, when appropriate, to avoid unnecessary iteration. The related theory and a description of the algorithms will be found in [8].

3.6 An E/R List Construction Algorithm

Numeric calculations in the global algorithm are driven by what we have called an E/R list. If there are n factors in the full analysis-of-variance model, then the E/R list will have 2^n entries corresponding to the 2^n terms (mean, main effects, and interactions) in a full factorial model. The E/R list is constructed by scanning or parsing an algebraic model statement. Numeric entries are made in the list which uniquely describe the model. Following the procedure suggested by Hemmerle [14], decreasing powers of 2 are assigned as numerical values to factor symbols and their associated subscripts, with unity for the last factor and its associated subscript.

The sum of the numerical values of the factor symbols (plus one) is computed for each term in the model. For crossed factors or interactions between crossed factors, this sum is entered in the corresponding location of the E/R list. For nested factors or interactions involving nested factors, multiple entries are made in the E/R list. The sum is entered into each location of the E/R list which corresponds to an array in the factorial decomposition which must be pooled (see [13, pp. 174-176]). The algorithm determines these locations from the numerical values of the factor symbols and subscripts in the model term. Examples will be given later in the sequel.

An hypothesis statement indicates those terms in the full model which should be deleted in forming the reduced model. As this statement is parsed, the E/R list entries for the model terms in the statement are made negative. The R operator is then always formed by pooling the arrays in the factorial decomposition associated with zero or negative E/R list entries.

4. STATISTICAL FOUNDATIONS AND COMPUTATIONS

Analysis-of-variance models tend to create confusion because of the fact that in the form they are usually written, these models are overparameterized. For example, the two-way classification with interaction is usually written as

$$y_{ijk} = \mu + a_i + b_j + ab_{ij} + e_{ijk},$$

$$i = 1, \dots, I, \quad j = 1, \dots, J, \quad k = 1, \dots, n_{ij}. \quad (4.1)$$

With e_{ijk} distributed $N(0, \sigma^2)$ and σ^2 unknown. None of the parameters μ , a_i , b_j , or ab_{ij} in (4.1) are estimable (for any i or j), even though the data may be balanced. That is, although the normal equations for (4.1) are consistent, there is no unique solution for these equations. We may obtain a unique solution by imposing constraints or side conditions upon the solution. These conditions have no effect upon what is estimable from the data for the model; however, they do have a bearing upon the hypothesis that is tested by applying the same applicable conditions to a reduced model.

Speed and Hocking [22] discuss the resulting difference in the application of Searle's $R(\cdot)$ notation [19] to the initial, overparameterized model as opposed to a reparameterized model obtained by imposing conditions on the original model. They refer to the latter as $R^*(\cdot)$ or procedure 2. Speed et al. [21] survey the hypotheses being tested by current computational methods and computer programs. A significant by-product of this work is that, when using procedure 2, the unweighted summation constraints (e.g., $\sum_i a_i = \sum_j b_j = \sum_i ab_{ij} = \sum_j ab_{ij} = 0$) yield the most likely hypotheses of interest, the classical Yates hypotheses [24].

The global algorithm uses procedure 2 with the unweighted summation conditions. In matrix form, the full reparameterized model is given by

$$y = X_0\beta + e \quad (4.2)$$

where y is a vector of N observations; X_0 is an $N \times p$ matrix of $-1, 0, 1$ indicator variables consistent with the unweighted summation conditions; β is a vector of p parameters; and the elements of e are distributed $N(0, \sigma^2)$.

Note that the p parameters in β will be linear combinations of the parameters appearing in the overparameterized model. As an example, consider the two-way classification *without* interaction. We have that

$$\beta' = (\mu + a. + b., a_1 - a., \dots, a_{I-1} - a., b_1 - b., \dots, b_{J-1} - b.) \quad (4.3)$$

with the dot notation denoting a mean. If X_0 has full rank, then all of the elements of β are linearly independent estimable functions. A fact that sometimes causes confusion is that $a_I - a.$ and $b_J - b.$ are also estimable in the above example; however, they are not linearly independent of the elements in β . We mention this fact here since the global algorithm yields a G-inverse solution for the initial overparameterized model as well as a G-inverse solution for (4.2). These solutions are the same, except for the inclusion of additional nonlinear independent elements in the former solution.

In order to consider hypothesis testing, we partition the design matrix X_0 for the full model as $[X_{01} | X_{02}]$, where X_{01} is $N \times k$, and write the model (4.2) as

$$y = X_{01}\beta_1 + X_{02}\beta_2 + e. \quad (4.4)$$

The global algorithm will attempt to test the hypothesis

$$H_0: \beta_2 = 0 \quad (4.5)$$

by fitting the reduced model

$$y = X_{01}\beta_1 + e. \quad (4.6)$$

The hypothesis statement specifies the model terms which are to be deleted from the full model in forming the reduced model (4.6). For balanced data or unbalanced data with no missing cells, these are the same hypotheses tested, with F statistics, in a standard balanced analysis-of-variance table. A rule is given in [9] to determine whether or not the equivalent balanced data hypothesis is testable for unbalanced data with missing cells. For a model with all factors crossed, this rule is simply: the equivalent balanced data hypothesis is testable if and only if the rank of the full model minus the rank of the reduced model equals the degrees of freedom one would use for the hypothesis with balanced data. It is the author's opinion that when the equivalent balanced data hypothesis is not testable, those hypotheses that are testable are usually very difficult to interpret in a meaningful way.

Some designs which are considered to be balanced designs will require iteration because the residual operator used applies to a more general model. An example of this is a Latin square which the algorithm treats as a missing cell problem. Iteration will be necessary to obtain SSE for the full model; however, as a result of data restructuring, iteration is not required in testing the relevant hypotheses.

The following statistics may be obtained from the algorithm, most of them on an optional basis:

- (1) cell sums, frequencies, and means;
- (2) classification sums, frequencies, and means;
- (3) rank (X_0);
- (4) SSE (full model) = $y'[I - X_0(X_0'X_0)^-X_0']y$ (invariant);
- (5) SSR (full model) = $y'X_0(X_0'X_0)^-X_0'y$;
- (6) a G-inverse solution to $X_0'X_0\hat{\beta} = X_0'y$; $\hat{\beta} = (X_0'X_0)^-X_0'y$;
- (7) estimates of expected cell means = $X_0\hat{\beta}$ (invariant);
- (8) rank (X_{01});

- (9) SSE (reduced model) = $y'[I - X_{01}(X'_{01}X_{01})^{-1}X'_{01}]y$;
 (10) SSR (reduced model) = $y'X_{01}(X'_{01}X_{01})^{-1}X'_{01}y$;
 (11) $F = \frac{(\text{SSR}(\text{full}) - \text{SSR}(\text{reduced})) / (\text{rank}(X_0) - \text{rank}(X_{01}))}{\text{SSE}(\text{full}) / (N - \text{rank}(X_0))}$;
 (12) probability values; $\text{Prob}\{F | H_0\}$.

We indicated earlier that a G-inverse solution is obtained for the overparameterized model as well. It should be clear that the rank of X_0 is the same as the rank of the design matrix of the overparameterized model formed using 0, 1 indicator variables.

5. LOGICAL COMPONENTS

The program for the global algorithm consists of a main program and eight subroutines. The principal function of these components is described below.

MAINA processes the factor, level statement; computes cell sums, cell frequencies, and $y'y$ as it reads the data; processes the options statement and contains much of the code to execute the A (cell means) and C (classification means) options.

SCAN processes the model and hypothesis statement to construct (or modify) the E/R list; computes parameters needed in restructuring data; computes the degrees of freedom applicable to data with no missing cells.

IGET is used by the main program and subroutine **SCAN** to sequentially retrieve characters (other than blank, plus, or comma) from the input buffer.

PART1 restructures the data (cell frequencies) when appropriate; checks for balance and alternative noniterative computations; computes rank noniteratively if possible or iteratively otherwise when the R (rank) option is specified.

PART2 is the principal numeric computational component; with the use of the remaining four subroutines, it computes SSE, SSR, estimates of expected cell means, a G-inverse solution to the normal equations, F statistics, and probability values.

STEP performs one basic step of the iterative AOV algorithm [12] using the improved approximation for SSR given in [10]. This basic step consists of the following substeps.

- (1) $A \leftarrow (Y - D \cdot V) / c$
- (2) $V \leftarrow V + A$
- (3) $B \leftarrow B + A$
- (4) $A \leftarrow R[A]$
- (5) $V \leftarrow V - A$
- (6) $S \leftarrow 2 \cdot Y'V - V'DV$

where

Y is the vector of cell sums, D is the diagonal matrix of cell frequencies stored as a vector, and c is an algorithm constant set to assure convergence and, in hypothesis testing, to assure monotonicity of the approximation S to SSR;

A , B , and V are work vectors of size n_c , with n_c being the number of cells. An additional work vector large enough to contain the factorial decomposition cited earlier is also required; however, vector A will occupy the first n_c locations of this additional vector;

$R[A]$ is the residual operator applied to vector A . Substep (4) above consists of a factorial decomposition of vector A , followed by pooling the appropriate arrays of the decomposition back into A .

The vectors B and V must be initially set to zero. Upon completion, perhaps after many steps, estimates of the expected cell means will be contained in V , and a G-inverse solution is obtained by applying the E operator to the vector B .

Substep (1) above is modified slightly for iterative rank computations and for restructured data. In computing rank, this takes the form

$$A \leftarrow (Y - \Delta V),$$

where Δ is a diagonal matrix with unit diagonals for filled cells and zero diagonals for missing cells. Operations with Δ are handled implicitly and this matrix is never stored. In performing noniterative calculations (a single step) on restructured cell frequencies, substep (1) is modified to

$$A \leftarrow (D2^- \cdot Y - V)$$

where $D2^-$ is a diagonal G-inverse of the diagonal matrix of restructured cell frequencies [8].

Substep (6) above must also be modified in computing rank iteratively with Δ replacing D . Furthermore, in rank computations, the vector Y contains a dummy unit vector; cell sums have been temporarily stored elsewhere in a vacant vector.

DECOMP obtains a factorial decomposition of a given vector; determines the classification frequencies needed in **PART1** to restructure data; computes classification means for the C option in **MAINA**.

POOL either moves the secondary array into the primary array, duplicating entries where needed, or it pools the secondary array into the primary array by addition.

LABEL calculates the array of coefficients for the array map needed in pooling; produces output labels for classification means and for the G-inverse solution.

Of these logical components, all of **PART1** may be deleted to further reduce program storage without making **PART2** inoperable; however, this would make the program less efficient owing to unnecessary iteration and full rank would be assumed by default for X_0 or X_{01} . Almost all of **SCAN** can also be deleted, provided the user constructs and inputs the required E/R list.

6. ARRAY STORAGE

We previously emphasized the storage efficiency of the algorithm for moderate to large problems. Fundamentally, this efficiency is attributable to the use of the iterative AOV algorithm; however, array storage economy was also stressed in constructing the global algorithm. The following is a brief description of all of the arrays included.

LSTFI is the array discussed in [13] and [18], which is used in formation and subsequent manipulation of the arrays in the factorial decomposition; this is an array of size 2^n , where n is the number of factors.

LER is the E/R list that is also of size 2^n .

LE is an alphanumeric array for factor symbols of size n .

LS is an alphanumeric array of associated subscripts of size n .

LV is an array of the numerical values assigned to the factor symbols and their associated subscripts; it also has size n .

LLIM contains the number of levels for each factor and is of size n .

LT is a temporary work array of size n .

LP is also a work array; however, its size is 10, the maximum number of factors handled by the program.

LD is an alphanumeric array of size 10 containing the digits 0-9; this array is used to convert alphanumeric input, such as the number of levels for a factor, to numeric.

LO is an alphanumeric array of size 10 containing the single-character option identifiers.

IA is an alphanumeric input buffer—a card image—which is used in processing the factor, levels, options, model, and hypothesis statements and also for storage

Table I. Options

Identifier	Function of option	Default
S(<i>d</i>)	Significant digits in results	<i>d</i> = 5
T(<i>l</i> × 10 ⁴)	Test level	<i>l</i> = 0.05
I(<i>m</i>)	Iteration maximum	<i>m</i> = 100
R	Rank computations (iterative)	off
V	Estimates of expected cell means	off
G	G-inverse solution to normal equations	off
P	Probability values for F statistics	off
Z	Intermediate output	off
A	Cell sums, frequencies, and means	off ^a
C	Classification sums, frequencies, and means	off ^a

^a These switches do not change their status each time the option is specified.

of a variable **FORMAT** statement for input data. (The model and hypothesis statements have continuation facilities.)

Q is the only two-dimensional array. It is used exclusively for rank computations and its size, in relation to the number of missing cells, determines the rank algorithm applied. If the dimension of **Q** were fixed at 10 × 10, then the quadratically convergent rank algorithm would be used only if there were no more than 10 missing cells.

QT is a vector whose size is the same as one of the dimensions of **Q**; it is used in conjunction with **Q** in computing rank.

W is a linear array that is used for all the numeric computations. This array logically consists of six contiguous vectors. The implicit names of these vectors coincide with those used in describing a step of the iterative AOV algorithm; they are ordered in storage as

Y, D, D2, V, B, A.

The first five of these vectors have size n_c , where n_c is the number of cells. The sixth vector is the size of the factorial decomposition for the problem. For example, with two factors we have $n_c = \mathbf{I} \cdot \mathbf{J}$, while the vector **A** has size $(\mathbf{I} + 1)(\mathbf{J} + 1)$. Ordinarily, the array **W** would be allocated all of the remaining computer storage; the implicit variable dimensioning of **W** maximizes the number of analysis-of-variance problems that can be processed within a given amount of storage.

7. THE FACTOR, LEVELS STATEMENT

The factor, levels statement creates entries for the arrays **LE**, **LS**, **LV**, and **LLIM**. The array **LSTFI** is then formed from **LLIM**. Factor symbols, subscript symbols, and factor levels are included in the statement; the number of factors and their numerical values are determined. As an example of a factor, levels statement, consider an analysis involving 3 factors with levels 5, 10, and 15. If we decide to name the factors **R**, **C**, **T** with subscripts **I**, **J**, **K**, then the statement would be written (typed/punched) on one line as

F(R, C, T), L(I(5), J(10), K(15))

All of the commas above are cosmetic, as are blanks and pluses; the same statement could be written as

F(RCT)L(I(5)J(10)K(15))

8. OPTIONS AND THE OPTIONS STATEMENT

A list of option identifiers and default values is given in Table I. The first three identifiers include an argument enclosed within parentheses which are used to modify a parameter. The remaining seven identifiers trigger on/off switches; with the exception of the A and C option, these switches change their status each time the option is specified.

The S option sets the relative iterative tolerance for SSR by forming the test constant $0.05 \times 10^{-d} \times (y'y)$. Iteration ceases when two successive approximations to SSR differ in magnitude by less than this test constant. This approach in practice seems to provide at least the requested accuracy in results obtained iteratively. A specification of $d = 0$ results in the skipping of all calculations in **PART2** and thus provides an option for independent rank calculations.

The T option sets the probability level at which F tests will be conducted. When iterative calculations are necessary for hypothesis testing, the constant c of the iterative AOV algorithm [12] is selected such that the sequence of approximations to the F statistic will be monotonically decreasing. The approximation given in [20] is used to compute the probability level for the F approximation. Whenever this value falls below the level specified, iteration ceases due to lack of significance. In the event that this value stays above the level specified, iteration continues until two successive F approximations differ by less than 0.005. When the P option is specified, iteration is continued for nonsignificant F's as well, in order to also obtain final probability levels for these statistics.

The R option must be in on status to compute rank *iteratively*. Noniterative rank computations in **PART1** are not affected by this option. The terminal iterative approximation to rank need not be highly accurate since the approximations converge monotonically to an integer; furthermore, absolute rather than relative cutoff tolerances may be used. Although the preset values for these tolerances may be modified, they have been tested to obtain the correct rank with minimal iteration.

The method best suited to a relatively small number of missing cells involves computing powers of a matrix whose order is the number of missing cells. The initial matrix is formed noniteratively. Successive powers, 1, 2, 4, 8, . . . , of this matrix are formed in the array **Q**. The trace of the matrix in **Q** converges quadratically and monotonically to the rank of X_0 . In this case, we cease iteration when successive traces differ by less than 0.1. The rank is taken as the next integer. The method best suited to a large number of missing cells applies the iterative AOV algorithm to dummy unit vectors for each filled cell with D replaced by Δ . The sum of the SSRs obtained for each unit vector will equal the rank of X_0 . In this case, we cease iteration when two successive approximations to SSR for a given unit vector differ by less than $0.1/n_f$ where n_f is the number of filled cells. This method will be expensive timewise when n_f is large.

When the Z option is on, the following intermediate output is produced when applicable: the current E/R list; the successive traces of the matrix in array **Q** or the cumulative sum of the SSRs for unit vectors in computing rank; the approximations to SSR for the full model; the approximations to SSR for the reduced model of the hypothesis.

Although the A and C options may appear to be somewhat mundane, these options require no additional array storage and they require little additional code; furthermore, it is likely that they would be requested as part of an analysis of variance.

An option statement is written on one line as a string of selected identifiers enclosed in parentheses and preceded by an O. The identifiers may be in any order. For example, with the R and Z options in off status, to specify eight significant digits of accuracy, iterative rank computations, and intermediate output, the options statement may be written as

O(R, S(8), Z)

where the commas are again cosmetic.

9. MODEL AND HYPOTHESIS STATEMENTS

Examples of the model and hypothesis statements for three factors are given in Table II along with the E/R list entries that would be produced for these statements. These examples assume that factor symbols A, B, C, and subscript symbols I, J, K were used in the factor, levels statement. The analysis-of-variance

Table II. Model and Hypothesis Statements

Statements	E/R list							
$(y_{ijkl} = \mu + a_i + b_j + c_k + e_{ijkl})$	ABC	AB	AC	A	BC	B	C	M
M + A(I) + B(J) + C(K)	0	0	0	5	0	3	2	1
H B(J)	0	0	0	5	0	-3	2	1
$(y_{ijkl} = \mu + a_i + b_{ij} + c_{ijk} + e_{ijkl})$	ABC	AB	AC	A	BC	B	C	M
M + A(I) + B(IJ) + C(IJK)	2	3	2	5	2	3	2	1
H C(IJK)	-2	3	-2	5	-2	3	-2	1
$(y_{ijkl} = \mu + a_i + b_{ij} + c_k + ac_{ik} + e_{ijkl})$	ABC	AB	AC	A	BC	B	C	M
M + A(I) + B(IJ) + C(K) + AC(IK)	0	3	6	5	0	3	2	1
H AC(IK)	0	3	-6	5	0	3	2	1
$(y_{ijkl} = \mu + a_i + b_j + ab_{ij} + c_{ijk} + e_{ijkl})$	ABC	AB	AC	A	BC	B	C	M
M + A(I) + B(J) + AB(IJ) + C(IJK)	2	7	2	5	2	3	2	1
H AB(IJ), C(IJK)	-2	-7	-2	5	-2	3	-2	1

model, as it is usually written, appears above each model statement in Table II. Notice that the variate y_{ijkl} and the error term e_{ijkl} do not appear in the model statement. The additional subscript l is also implicit; that is, the number of observations in the (i, j, k) cell depends upon the data. The rules for constructing the model statement basically parallel those used in [14] and [18]: the associated subscript of a factor must always appear along with the factor symbol in a model term; if a factor is nested, it must be nested within a factor or factors appearing to the left of it in the factor, levels statement; the ordering of the associated subscripts in the factor, levels statement must correspond with the ordering of the data. As indicated earlier, pluses, blanks, and commas serve only to improve the appearance of these statements; these characters are ignored in parsing. Should either the model or hypothesis statement require more than one line (or card), continuation to the next line is indicated by placing a slash (/) after the last model term in the current line. For example, the fourth model statement in Table II could be written as

$$\mathbf{M + A(I) + B(J) + AB(IJ) / + C(IJK)}$$

As a convenience, a full factorial model may be specified as M^* , an asterisk following M . Hypothesis statements are then written normally using the same symbols appearing in the factor, levels statement. The E/R list construction algorithm includes the logic necessary to detect invalid models or hypotheses incorrectly specified. During interactive use, corrections can be made when an error is detected.

Any number of models may be applied experimentally to a given set of data. This includes models with fewer factors than appear in the factor, levels statement; the data will automatically be restructured for such models.

10. FLOW OF CONTROL

Flow of control proceeds in the following manner:

- (1) factor, levels statement,
- (2) variable **FORMAT** statement,
- (3) data,
- (4) blank line (card),
- (5) options statement/model statement/hypothesis statement/blank line/E.

The data are prepared with one observation per line and cell indicators preceding the observation; the variable **FORMAT** statement must reflect this order. Cell

indicators are the factor levels for the observation. There is no indicator for within cell replication and no entries are made for missing cells. A check on lexicographical ordering of indicators is included; however, this check is easily deleted if inconvenient.

Once the data have been entered, control is centered upon (5). Any number of options, model, or hypothesis statements may be entered in any order (except that a model statement must precede any hypothesis statements). Entering a blank line returns control to (1); entering the character E, for end, will terminate processing.

11. EXAMPLE

We present here a simple example, processed in batch mode, to illustrate many of the features of the global algorithm. There are three factors A, B, and C with levels of 3, 2, and 2, respectively; this information is conveyed in the initial factor, levels statement. The data are unbalanced and three cells—cells (1, 1, 2), (2, 1, 2), and (2, 2, 1)—are empty.

After the data have been read, the first option statement is entered requesting intermediate output and specifying the A, C, and R options. Summary statistics, computed immediately for the A and C options, are printed followed by the current status of the options, including default values. A model statement representing a model with three crossed factors A, B, and C and an interaction between factors A and B is entered; the corresponding E/R list is printed as intermediate output. Intermediate output of rank computations for the design matrix for the full model follows, with each iteration displaying the trace of the matrix in **Q**, until the rank of 7 has been determined. SSR for the model is next computed iteratively to approximately five significant digits with intermediate output printed for each iteration.

An hypothesis statement is entered to test the significance of factor C. The algorithm determines the rank of the design matrix for the reduced model to be six noniteratively; it also restructures the data and computes SSR noniteratively for the reduced model. The computed F statistic is found to be not significant at the 5 percent level. Another hypothesis statement is entered to test the hypothesis of no AB interaction. Again the algorithm determines rank noniteratively; however, it must iterate to determine SSR for the reduced model. Intermediate output of the monotonically increasing approximation to SSR is printed until the corresponding approximation for F falls below the 5 percent level on the second iteration.

```
F(A,B,C) L(I(3),J(2),K(2))
DATA FORMAT AND INPUT DATA-
(1X,3I2,F4.1)
1 1 1 2.2
1 1 1 2.8
1 2 1 3.1
1 2 2 4.5
2 1 1 2.4
2 1 1 2.7
2 2 2 2.6
2 2 2 5.2
3 1 1 4.5
3 1 2 3.6
3 2 1 5.0
3 2 2 3.7
0 0 0 .0
O(Z,A,C,R)
CELL SUMS, FREQUENCIES, AND MEANS-
CELL      SUM      FREQ.      MEAN
1          .50000000e+01  2.          .25000000e+01
2          (MISSING CELL)
3          .31000000e+01  1.          .31000000e+01
4          .45000000e+01  1.          .45000000e+01
5          .51000000e+01  2.          .25500000e+01
6          (MISSING CELL)
7          (MISSING CELL)
8          .78000000e+01  2.          .39000000e+01
9          .45000000e+01  1.          .45000000e+01
10         .36000000e+01  1.          .36000000e+01
11         .50000000e+01  1.          .50000000e+01
12         .37000000e+01  1.          .37000000e+01
```

CLASSIFICATION SUMS, FREQUENCIES, AND MEANS-
I.J.

1	.50000000e+01	2.	.25000000e+01
2	.76000000e+01	2.	.38000000e+01
3	.51000000e+01	2.	.25500000e+01
4	.78000000e+01	2.	.39000000e+01
5	.81000000e+01	2.	.40500000e+01
6	.87000000e+01	2.	.43500000e+01

I.K

1	.81000000e+01	3.	.27000000e+01
2	.45000000e+01	1.	.45000000e+01
3	.51000000e+01	2.	.25500000e+01
4	.78000000e+01	2.	.39000000e+01
5	.95000000e+01	2.	.47500000e+01
6	.73000000e+01	2.	.36500000e+01

I..

1	.12600000e+02	4.	.31500000e+01
2	.12900000e+02	4.	.32250000e+01
3	.16800000e+02	4.	.42000000e+01

.JK

1	.14600000e+02	5.	.29200000e+01
2	.36000000e+01	1.	.36000000e+01
3	.81000000e+01	2.	.40500000e+01
4	.16000000e+02	4.	.40000000e+01

.J.

1	.18200000e+02	6.	.30333333e+01
2	.24100000e+02	6.	.40166667e+01

..K

1	.22700000e+02	7.	.32428571e+01
2	.19600000e+02	5.	.39200000e+01

...

1	.42300000e+02	12.	.35250000e+01
---	---------------	-----	---------------

OPTIONS- S= 5, T= .0500, I=100, R=1, V=0, G=0, P=0

M+A(I)+B(J)+AB(IJ)+C(K)

E/R LIST-

0	7	0	5	0	3	2	1
---	---	---	---	---	---	---	---

ITERATION 0, TRACE= 5.25000000

ITERATION 1, TRACE= 5.93750000

ITERATION 2, TRACE= 6.558593750

ITERATION 3, TRACE= 6.892074585

THE RANK OF THE M DESIGN MATRIX IS 7

ITERATION 1, SSRM= .11689479e+03

ITERATION 2, SSRM= .14589862e+03

ITERATION 3, SSRM= .15314517e+03

ITERATION 4, SSRM= .15495693e+03

ITERATION 5, SSRM= .15541004e+03

ITERATION 6, SSRM= .15552343e+03

ITERATION 7, SSRM= .15555187e+03

ITERATION 8, SSRM= .15555904e+03

ITERATION 9, SSRM= .15556089e+03

ITERATION 10, SSRM= .15556138e+03

ITERATION 11, SSRM= .15556153e+03

ITERATION 12, SSRM= .15556158e+03

ITERATION 12, SSR(FULL MODEL)= .15556158e+03,

SSE(FULL MODEL)= .57284167e+01

H C(K)

E/R LIST-

0	7	0	5	0	3	-2	1
---	---	---	---	---	---	----	---

THE RANK OF THE H DESIGN MATRIX IS 6

FROM RANK COMPUTATIONS- DF(NUM)= 1, DF(DEN)= 5

ITERATION 1, SSRH= .15545500e+03

ITERATION 1, F= .093, PROB(F) .GT. .7633 VS. F LEVEL OF .0500

SSR(REDUCED MODEL)= .15545500e+03

H AB(IJ)

E/R LIST-

0	-7	0	5	0	3	2	1
---	----	---	---	---	---	---	---

THE RANK OF THE H DESIGN MATRIX IS 5

FROM RANK COMPUTATIONS- DF(NUM)= 2, DF(DEN)= 5

ITERATION 1, SSRH= .11623292e+03

ITERATION 2, SSRH= .14498170e+03

ITERATION 2, F= 4.617, PROB(F) .GT. .0734 VS. F LEVEL OF .0500

SSR(REDUCED MODEL)= .14498170e+03

O(Z,V,G)

OPTIONS- S= 5, T= .0500, I=100, R=1, V=1, G=1, P=0

M+A(I)+B(J)

THE RANK OF THE M DESIGN MATRIX IS 4

ITERATION 1, SSR(FULL MODEL)= .15475333e+03,

SSE(FULL MODEL)= .65366667e+01

ESTIMATES OF EXPECTED CELL MEANS-

CELL	ESTIMATED MEAN
1	.26583333e+01
2	.26583333e+01 (MISSING CELL)
3	.36416667e+01
4	.36416667e+01
5	.27333333e+01
6	.27333333e+01 (MISSING CELL)
7	.37166667e+01 (MISSING CELL)
8	.37166667e+01
9	.37083333e+01

```

10 .37083333e+01
11 .46916667e+01
12 .46916667e+01
G-INVERSE SOLUTION
A
  1 -.37500000e+00
  2 -.30000000e+00
  3 .67500000e+00
B
  1 -.49166667e+00
  2 .49166667e+00
E
  1 .35250000e+00

```

Another options statement is entered which suspends intermediate output and specifies the V and G option; again the status of the options is printed. A model statement representing the two-way classification without interaction for factors A and B is entered. All of the remaining computations are noniterative. Estimates of expected cell means are printed followed by a solution to the normal equations in response to the V and G options. Unfortunately, this example does not convey the algorithm's facility to handle large data sets and diverse models.

12. SUMMARY

We have described a comprehensive algorithm for analysis of variance which has the facility to handle large problems involving unbalanced data on small computers. At the same time, the computational efficiencies common to balanced data are realized. Array storage economy is achieved by virtue of the fact that the algorithm performs vector operations as opposed to matrix operations. Neither the design matrix of indicator variables X_0 nor the matrix $X_0'X_0$ is ever formed or stored. Ancillary computations and output amenities have been deliberately limited to conserve program storage; however, the facility for algebraic problem specifications has been included.

For some data sets and models requiring iteration, the algorithm may converge slowly so that it will be expensive timewise to achieve a high degree of accuracy; however, performance studies of the iterative AOV algorithm given in [10] and [6] indicate that the algorithm is likely to outperform matrix-oriented algorithms timewise on large problems with unbalanced data. It should always outperform these algorithms timewise when the data are balanced. Subsequent to submission of this algorithm, other authors [2], [15] have suggested using the conjugate gradient method rather than iterative improvement for the basic iterative step. The monotonicity properties of the algorithm are preserved with the conjugate gradient method, and performance should be improved whenever iteration for SSR is necessary. Subroutine step may be modified relatively easily to reflect using this method.

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ALGORITHM

[A part of the listing is printed here. The complete listing is available from the ACM Distribution Service.]

C	**A COMPREHENSIVE, MATRIX FREE ALGORITHM FOR ANALYSIS OF VARIANCE**	MAN	10
C		MAN	20
C	**STATISTICS COMPUTED**	MAN	30
C		MAN	40
C	1) CELL SUMS, FREQUENCIES, AND MEANS	MAN	50
C		MAN	60
C	2) CLASSIFICATION SUMS, FREQUENCIES, AND MEANS	MAN	70
C		MAN	80
C	3) RANK(DESIGN MATRIX FOR FULL MODEL)	MAN	90
C		MAN	100
C	4) SSE(FULL MODEL) AND SSR(FULL MODEL)	MAN	110
C		MAN	120
C	5) A G-INVERSE SOLUTION TO THE NORMAL EQUATIONS	MAN	130
C		MAN	140
C	6) ESTIMATES OF EXPECTED CELL MEANS	MAN	150
C		MAN	160
C	7) RANK(RESTRICTED DESIGN MATRIX FOR REDUCED MODEL)	MAN	170
C		MAN	180
C	8) SSE(REDUCED MODEL) AND SSR(REDUCED MODEL)	MAN	190
C		MAN	200
C	9) THE F STATISTIC FOR A SPECIFIED HYPOTHESIS	MAN	210

C		MAN	220
C	10) PROBABILITY OF A GREATER F GIVEN THE HYPOTHESIS	MAN	230
C		MAN	240
C	**FLOW OF CONTROL**	MAN	250
C		MAN	260
C	1) FACTOR, LEVELS STATEMENT	MAN	270
C		MAN	280
C	2) VARIABLE FORMAT STATEMENT	MAN	290
C		MAN	300
C	3) DATA	MAN	310
C		MAN	320
C	4) BLANK LINE (CARD)	MAN	330
C		MAN	340
C	5) OPTIONS STATEMENT/MODEL STATEMENT/HYPOTHESIS STATEMENT/BLANK	MAN	350
C	LINE/E	MAN	360
C		MAN	370
C	E ENDS PROCESSING, BLANK LINE SENDS CONTROL TO 1), STATEMENTS	MAN	380
C	RETURN CONTROL TO 5) FOLLOWING EXECUTION.	MAN	390
C		MAN	400
C	**FACTOR, LEVELS STATEMENT**	MAN	410
C		MAN	420
C	A DISTINCT ALPHABETIC CHARACTER IS USED TO NAME EACH FACTOR AND	MAN	430
C	EACH SUBSCRIPT. THE STRING OF FACTOR SYMBOLS IS PLACED IN PARENS	MAN	440
C	AND THEN PRECEDED BY THE LETTER F. THIS IS FOLLOWED BY THE LET-	MAN	450
C	TER L AND A PARENTHEZIZED LIST OF ASSOCIATED SUBSCRIPTS WITH THE	MAN	460
C	NUMBER OF LEVELS FOR THE FACTOR IN PARENS FOLLOWING THE ASSOCI-	MAN	470
C	ATED SUBSCRIPT.	MAN	480
C		MAN	490
C	EXAMPLE FOR 3 FACTORS NAMED A, B, C, ASSOCIATED SUBSCRIPTS NAMED	MAN	500
C	I, J, K, AND NUMBER OF LEVELS OF THE FACTORS EQUAL TO 5, 10, 15:	MAN	510
C		MAN	520
C	F(A,B,C) L(I(5),J(10),K(15))	MAN	530
C		MAN	540
C	COMMAS AND BLANKS ARE COSMETIC	MAN	550
C		MAN	560
C	ORDERING OF THE ASSOCIATED SUBSCRIPTS IS ASSUMED TO CORRESPOND	MAN	570
C	TO THE ORDERING OF THE DATA. THE NUMBER OF OBSERVATIONS WITHIN	MAN	580
C	THE (I,J,K) CELL DEPENDS UPON THE DATA; NO SUBSCRIPT SYMBOL IS	MAN	590
C	USED TO DENOTE WITHIN CELL REPLICATION.	MAN	600
C		MAN	610