

Matrix

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theory MatrixGeneral imports Main begin

types 'a infmatrix = [nat, nat]  $\Rightarrow$  'a

constdefs
  nonzero-positions :: ('a::zero) infmatrix  $\Rightarrow$  (nat*nat) set
  nonzero-positions A == {pos. A (fst pos) (snd pos)  $\sim$  0}

typedef 'a matrix = {(f::('a::zero) infmatrix). finite (nonzero-positions f)}
apply (rule-tac x=( $\% j i. 0$ ) in exI)
by (simp add: nonzero-positions-def)

declare Rep-matrix-inverse[simp]

lemma finite-nonzero-positions : finite (nonzero-positions (Rep-matrix A))
apply (rule Abs-matrix-induct)
by (simp add: Abs-matrix-inverse matrix-def)

constdefs
  nrows :: ('a::zero) matrix  $\Rightarrow$  nat
  nrows A == if nonzero-positions(Rep-matrix A) = {} then 0 else Suc(Max
((image fst) (nonzero-positions (Rep-matrix A))))
  ncols :: ('a::zero) matrix  $\Rightarrow$  nat
  ncols A == if nonzero-positions(Rep-matrix A) = {} then 0 else Suc(Max ((image
snd) (nonzero-positions (Rep-matrix A))))

lemma nrows:
  assumes hyp: nrows A  $\leq$  m
  shows (Rep-matrix A m n) = 0 (is ?concl)
proof cases
  assume nonzero-positions(Rep-matrix A) = {}
  then show (Rep-matrix A m n) = 0 by (simp add: nonzero-positions-def)
next
  assume a: nonzero-positions(Rep-matrix A)  $\neq$  {}
  let ?S = fst'(nonzero-positions(Rep-matrix A))
  from a have b: ?S  $\neq$  {} by (simp)
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have  $c$ : finite (? $S$ ) by (simp add: finite-nonzero-positions)
from hyp have  $d$ : Max (? $S$ ) <  $m$  by (simp add: a n rows-def)
have  $m \notin ?S$ 
proof –
  have  $m \in ?S \implies m \leq \text{Max}(?S)$  by (simp add: Max-ge[OF c b])
  moreover from  $d$  have  $\sim(m \leq \text{Max } ?S)$  by (simp)
  ultimately show  $m \notin ?S$  by (auto)
qed
thus Rep-matrix  $A$   $m$   $n = 0$  by (simp add: nonzero-positions-def image-Collect)
qed

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constdefs

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transpose-infmatrix :: 'a infmatrix  $\Rightarrow$  'a infmatrix
transpose-infmatrix  $A$   $j$   $i == A$   $i$   $j$ 
transpose-matrix :: ('a::zero) matrix  $\Rightarrow$  'a matrix
transpose-matrix == Abs-matrix o transpose-infmatrix o Rep-matrix

```

declare *transpose-infmatrix-def*[*simp*]

lemma *transpose-infmatrix-twice*[*simp*]: *transpose-infmatrix* (*transpose-infmatrix* A) = A
by (*(rule ext)+, simp*)

lemma *transpose-infmatrix*: *transpose-infmatrix* (% j i . P j i) = (% j i . P i j)
apply (*rule ext*)
by (*simp add: transpose-infmatrix-def*)

lemma *transpose-infmatrix-closed*[*simp*]: *Rep-matrix* (*Abs-matrix* (*transpose-infmatrix* (*Rep-matrix* x))) = *transpose-infmatrix* (*Rep-matrix* x)

apply (*rule Abs-matrix-inverse*)
apply (*simp add: matrix-def nonzero-positions-def image-def*)

proof –

let ? $A = \{pos. \text{Rep-matrix } x \text{ (snd pos) (fst pos)} \neq 0\}$

let ?*swap* = % pos . (snd pos , fst pos)

let ? $B = \{pos. \text{Rep-matrix } x \text{ (fst pos) (snd pos)} \neq 0\}$

have *swap-image*: ?*swap*'? A = ? B

apply (*simp add: image-def*)

apply (*rule set-ext*)

apply (*simp*)

proof

fix y

assume *hyp*: $\exists a$ b . *Rep-matrix* x b $a \neq 0 \wedge y = (b, a)$

thus *Rep-matrix* x (fst y) (snd y) $\neq 0$

proof –

from *hyp* **obtain** a b **where** (*Rep-matrix* x b $a \neq 0 \ \& \ y = (b, a)$) **by** *blast*

then show *Rep-matrix* x (fst y) (snd y) $\neq 0$ **by** (*simp*)

qed

next

fix y

assume *hyp*: *Rep-matrix* x (*fst* y) (*snd* y) $\neq 0$
show $\exists a b. (\text{Rep-matrix } x \ b \ a \neq 0 \ \& \ y = (b,a))$ **by** (*rule* *exI*[*of* - *snd* y],
rule *exI*[*of* - *fst* y], *simp*)
qed
then have *finite* (*?swap*' $?A$)
proof –
have *finite* (*nonzero-positions* (*Rep-matrix* x)) **by** (*simp* *add*: *finite-nonzero-positions*)
then have *finite* $?B$ **by** (*simp* *add*: *nonzero-positions-def*)
with *swap-image* **show** *finite* (*?swap*' $?A$) **by** (*simp*)
qed
moreover
have *inj-on* *?swap* $?A$ **by** (*simp* *add*: *inj-on-def*)
ultimately show *finite* $?A$ **by** (*rule* *finite-imageD*[*of* *?swap* $?A$])
qed

lemma *infmatrixforward*: $(x::'a \text{ infmatrix}) = y \implies \forall a b. x \ a \ b = y \ a \ b$ **by** *auto*

lemma *transpose-infmatrix-inject*: $(\text{transpose-infmatrix } A = \text{transpose-infmatrix } B) = (A = B)$
apply (*auto*)
apply (*rule* *ext*)
apply (*simp* *add*: *transpose-infmatrix*)
apply (*drule* *infmatrixforward*)
apply (*simp*)
done

lemma *transpose-matrix-inject*: $(\text{transpose-matrix } A = \text{transpose-matrix } B) = (A = B)$
apply (*simp* *add*: *transpose-matrix-def*)
apply (*subst* *Rep-matrix-inject*[*THEN* *sym*])
apply (*simp* *only*: *transpose-infmatrix-closed* *transpose-infmatrix-inject*)
done

lemma *transpose-matrix[simp]*: *Rep-matrix*(*transpose-matrix* A) $j \ i = \text{Rep-matrix } A \ i \ j$
by (*simp* *add*: *transpose-matrix-def*)

lemma *transpose-transpose-id[simp]*: *transpose-matrix* (*transpose-matrix* A) = A
by (*simp* *add*: *transpose-matrix-def*)

lemma *nrows-transpose[simp]*: *nrows* (*transpose-matrix* A) = *ncols* A
by (*simp* *add*: *nrows-def* *ncols-def* *nonzero-positions-def* *transpose-matrix-def* *image-def*)

lemma *ncols-transpose[simp]*: *ncols* (*transpose-matrix* A) = *nrows* A
by (*simp* *add*: *nrows-def* *ncols-def* *nonzero-positions-def* *transpose-matrix-def* *image-def*)

lemma *ncols*: *ncols* $A \leq n \implies \text{Rep-matrix } A \ m \ n = 0$

proof –
assume *ncols* $A \leq n$

then have $nrows$ (*transpose-matrix A*) $\leq n$ **by** (*simp*)
then have *Rep-matrix* (*transpose-matrix A*) $n\ m = 0$ **by** (*rule nrows*)
thus *Rep-matrix A* $m\ n = 0$ **by** (*simp add: transpose-matrix-def*)
qed

lemma *ncols-le*: $(ncols\ A \leq n) = (!\ j\ i.\ n \leq i \longrightarrow (Rep\ matrix\ A\ j\ i) = 0)$ (**is** $- = ?st$)
apply (*auto*)
apply (*simp add: ncols*)
proof (*simp add: ncols-def, auto*)
let $?P = nonzero\ positions$ (*Rep-matrix A*)
let $?p = snd^i ?P$
have $a: finite\ ?p$ **by** (*simp add: finite-nonzero-positions*)
let $?m = Max\ ?p$
assume $\sim(Suc\ (?m) \leq n)$
then have $b: n \leq ?m$ **by** (*simp*)
fix $a\ b$
assume $(a, b) \in ?P$
then have $?p \neq \{\}$ **by** (*auto*)
with a **have** $?m \in ?p$ **by** (*simp*)
moreover have $!x.\ (x \in ?p \longrightarrow (?y.\ (Rep\ matrix\ A\ y\ x) \neq 0))$ **by** (*simp add: nonzero-positions-def image-def*)
ultimately have $?y.\ (Rep\ matrix\ A\ y\ ?m) \neq 0$ **by** (*simp*)
moreover assume $?st$
ultimately show *False* **using** b **by** (*simp*)
qed

lemma *less-ncols*: $(n < ncols\ A) = (?j\ i.\ n \leq i \ \&\ (Rep\ matrix\ A\ j\ i) \neq 0)$ (**is** $?concl$)
proof $-$
have $a: !! (a::nat)\ b.\ (a < b) = (\sim(b \leq a))$ **by** *arith*
show $?concl$ **by** (*simp add: a ncols-le*)
qed

lemma *le-ncols*: $(n \leq ncols\ A) = (\forall\ m.\ (\forall\ j\ i.\ m \leq i \longrightarrow (Rep\ matrix\ A\ j\ i) = 0) \longrightarrow n \leq m)$ (**is** $?concl$)
apply (*auto*)
apply (*subgoal-tac ncols A <= m*)
apply (*simp*)
apply (*simp add: ncols-le*)
apply (*drule-tac x=ncols A in spec*)
by (*simp add: ncols*)

lemma *nrows-le*: $(nrows\ A \leq n) = (!\ j\ i.\ n \leq j \longrightarrow (Rep\ matrix\ A\ j\ i) = 0)$ (**is** $?s$)
proof $-$
have $(nrows\ A \leq n) = (ncols\ (transpose\ matrix\ A) \leq n)$ **by** (*simp*)
also have $\dots = (!\ j\ i.\ n \leq i \longrightarrow (Rep\ matrix\ (transpose\ matrix\ A)\ j\ i) = 0)$
by (*rule ncols-le*)

also have $\dots = (! j i. n <= i \longrightarrow (\text{Rep-matrix } A \ i \ j) = 0)$ **by** (*simp*)
finally show $(\text{nrows } A <= n) = (! j i. n <= j \longrightarrow (\text{Rep-matrix } A \ j \ i) = 0)$ **by**
(*auto*)
qed

lemma *less-nrows*: $(m < \text{nrows } A) = (? j i. m <= j \ \& \ (\text{Rep-matrix } A \ j \ i) \neq 0)$
(is ?concl)
proof –
have $a: !! (a::\text{nat}) \ b. (a < b) = (\sim(b <= a))$ **by** *arith*
show *?concl* **by** (*simp add: nrows-le*)
qed

lemma *le-nrows*: $(n <= \text{nrows } A) = (\forall m. (\forall j i. m <= j \longrightarrow (\text{Rep-matrix } A \ j \ i) = 0) \longrightarrow n <= m)$ **(is ?concl)**
apply (*auto*)
apply (*subgoal-tac nrows A <= m*)
apply (*simp*)
apply (*simp add: nrows-le*)
apply (*drule-tac x=nrows A in spec*)
by (*simp add: nrows*)

lemma *nrows-notzero*: $\text{Rep-matrix } A \ m \ n \neq 0 \implies m < \text{nrows } A$
apply (*case-tac nrows A <= m*)
apply (*simp-all add: nrows*)
done

lemma *ncols-notzero*: $\text{Rep-matrix } A \ m \ n \neq 0 \implies n < \text{ncols } A$
apply (*case-tac ncols A <= n*)
apply (*simp-all add: ncols*)
done

lemma *finite-natarray1*: $\text{finite } \{x. x < (n::\text{nat})\}$
apply (*induct n*)
apply (*simp*)
proof –
fix n
have $\{x. x < \text{Suc } n\} = \text{insert } n \ \{x. x < n\}$ **by** (*rule set-ext, simp, arith*)
moreover assume $\text{finite } \{x. x < n\}$
ultimately show $\text{finite } \{x. x < \text{Suc } n\}$ **by** (*simp*)
qed

lemma *finite-natarray2*: $\text{finite } \{\text{pos. } (\text{fst pos}) < (m::\text{nat}) \ \& \ (\text{snd pos}) < (n::\text{nat})\}$
apply (*induct m*)
apply (*simp+*)
proof –
fix $m::\text{nat}$
let $?s0 = \{\text{pos. } \text{fst pos} < m \ \& \ \text{snd pos} < n\}$
let $?s1 = \{\text{pos. } \text{fst pos} < (\text{Suc } m) \ \& \ \text{snd pos} < n\}$
let $?sd = \{\text{pos. } \text{fst pos} = m \ \& \ \text{snd pos} < n\}$

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assume f0: finite ?s0
have f1: finite ?sd
proof –
  let ?f = % x. (m, x)
  have {pos. fst pos = m & snd pos < n} = ?f ‘ {x. x < n} by (rule set-ext,
simp add: image-def, auto)
  moreover have finite {x. x < n} by (simp add: finite-natarray1)
  ultimately show finite {pos. fst pos = m & snd pos < n} by (simp)
qed
have su: ?s0 ∪ ?sd = ?s1 by (rule set-ext, simp, arith)
from f0 f1 have finite (?s0 ∪ ?sd) by (rule finite-UnI)
with su show finite ?s1 by (simp)
qed

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lemma RepAbs-matrix:

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assumes aem: ? m. ! j i. m <= j → x j i = 0 (is ?em) and aen: ? n. ! j i. (n
<= i → x j i = 0) (is ?en)
shows (Rep-matrix (Abs-matrix x)) = x
apply (rule Abs-matrix-inverse)
apply (simp add: matrix-def nonzero-positions-def)
proof –
  from aem obtain m where a: ! j i. m <= j → x j i = 0 by (blast)
  from aen obtain n where b: ! j i. n <= i → x j i = 0 by (blast)
  let ?u = {pos. x (fst pos) (snd pos) ≠ 0}
  let ?v = {pos. fst pos < m & snd pos < n}
  have c: !! (m::nat) a. ~ (m <= a) ⇒ a < m by (arith)
  from a b have (?u ∩ (-?v)) = {}
    apply (simp)
    apply (rule set-ext)
    apply (simp)
    apply auto
  by (rule c, auto)+
  then have d: ?u ⊆ ?v by blast
  moreover have finite ?v by (simp add: finite-natarray2)
  ultimately show finite ?u by (rule finite-subset)
qed

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constdefs

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apply-infmatrix :: ('a ⇒ 'b) ⇒ 'a infmatrix ⇒ 'b infmatrix
apply-infmatrix f == % A. (% j i. f (A j i))
apply-matrix :: ('a ⇒ 'b) ⇒ ('a::zero) matrix ⇒ ('b::zero) matrix
apply-matrix f == % A. Abs-matrix (apply-infmatrix f (Rep-matrix A))
combine-infmatrix :: ('a ⇒ 'b ⇒ 'c) ⇒ 'a infmatrix ⇒ 'b infmatrix ⇒ 'c infmatrix
combine-infmatrix f == % A B. (% j i. f (A j i) (B j i))
combine-matrix :: ('a ⇒ 'b ⇒ 'c) ⇒ ('a::zero) matrix ⇒ ('b::zero) matrix ⇒
('c::zero) matrix
combine-matrix f == % A B. Abs-matrix (combine-infmatrix f (Rep-matrix A)
(Rep-matrix B))

```

lemma *expand-apply-infmatrix*[simp]: *apply-infmatrix* f A j i = f (A j i)
by (*simp add: apply-infmatrix-def*)

lemma *expand-combine-infmatrix*[simp]: *combine-infmatrix* f A B j i = f (A j i)
(B j i)
by (*simp add: combine-infmatrix-def*)

constdefs

commutative :: ($'a \Rightarrow 'a \Rightarrow 'b$) \Rightarrow *bool*
commutative f == ! x y . f x y = f y x
associative :: ($'a \Rightarrow 'a \Rightarrow 'a$) \Rightarrow *bool*
associative f == ! x y z . f (f x y) z = f x (f y z)

To reason about associativity and commutativity of operations on matrices, let's take a step back and look at the general situation: Assume that we have sets A and B with $B \subset A$ and an abstraction $u : A \rightarrow B$. This abstraction has to fulfill $u(b) = b$ for all $b \in B$, but is arbitrary otherwise. Each function $f : A \times A \rightarrow A$ now induces a function $f' : B \times B \rightarrow B$ by $f' = u \circ f$. It is obvious that commutativity of f implies commutativity of f' : $f'xy = u(fxy) = u(fyx) = f'yx$.

lemma *combine-infmatrix-commute*:
commutative $f \implies$ *commutative* (*combine-infmatrix* f)
by (*simp add: commutative-def combine-infmatrix-def*)

lemma *combine-matrix-commute*:
commutative $f \implies$ *commutative* (*combine-matrix* f)
by (*simp add: combine-matrix-def commutative-def combine-infmatrix-def*)

On the contrary, given an associative function f we cannot expect f' to be associative. A counterexample is given by $A = \mathbb{Z}$, $B = \{-1, 0, 1\}$, as f we take addition on \mathbb{Z} , which is clearly associative. The abstraction is given by $u(a) = 0$ for $a \notin B$. Then we have

$$f'(f'11) - 1 = u(f(u(f11)) - 1) = u(f(u2) - 1) = u(f0 - 1) = -1,$$

but on the other hand we have

$$f'1(f'1 - 1) = u(f1(u(f1 - 1))) = u(f10) = 1.$$

A way out of this problem is to assume that $f(A \times A) \subset A$ holds, and this is what we are going to do:

lemma *nonzero-positions-combine-infmatrix*[simp]: f 0 $0 = 0 \implies$ *nonzero-positions* (*combine-infmatrix* f A B) \subseteq (*nonzero-positions* A) \cup (*nonzero-positions* B)
by (*rule subsetI, simp add: nonzero-positions-def combine-infmatrix-def, auto*)

lemma *finite-nonzero-positions-Rep*[simp]: *finite* (*nonzero-positions* (*Rep-matrix* A))
by (*insert Rep-matrix [of A], simp add: matrix-def*)

lemma *combine-infmatrix-closed* [simp]:
 $f \ 0 \ 0 = 0 \implies \text{Rep-matrix } (\text{Abs-matrix } (\text{combine-infmatrix } f \ (\text{Rep-matrix } A) \ (\text{Rep-matrix } B))) = \text{combine-infmatrix } f \ (\text{Rep-matrix } A) \ (\text{Rep-matrix } B)$
apply (rule *Abs-matrix-inverse*)
apply (simp add: *matrix-def*)
apply (rule *finite-subset*[of - (*nonzero-positions* (*Rep-matrix* A)) \cup (*nonzero-positions* (*Rep-matrix* B))])
by (simp-all)

We need the next two lemmas only later, but it is analog to the above one, so we prove them now:

lemma *nonzero-positions-apply-infmatrix*[simp]: $f \ 0 = 0 \implies \text{nonzero-positions } (\text{apply-infmatrix } f \ A) \subseteq \text{nonzero-positions } A$
by (rule *subsetI*, simp add: *nonzero-positions-def* *apply-infmatrix-def*, auto)

lemma *apply-infmatrix-closed* [simp]:
 $f \ 0 = 0 \implies \text{Rep-matrix } (\text{Abs-matrix } (\text{apply-infmatrix } f \ (\text{Rep-matrix } A))) = \text{apply-infmatrix } f \ (\text{Rep-matrix } A)$
apply (rule *Abs-matrix-inverse*)
apply (simp add: *matrix-def*)
apply (rule *finite-subset*[of - *nonzero-positions* (*Rep-matrix* A)])
by (simp-all)

lemma *combine-infmatrix-assoc*[simp]: $f \ 0 \ 0 = 0 \implies \text{associative } f \implies \text{associative } (\text{combine-infmatrix } f)$
by (simp add: *associative-def* *combine-infmatrix-def*)

lemma *comb*: $f = g \implies x = y \implies f \ x = g \ y$
by (auto)

lemma *combine-matrix-assoc*: $f \ 0 \ 0 = 0 \implies \text{associative } f \implies \text{associative } (\text{combine-matrix } f)$
apply (simp(*no-asm*) add: *associative-def* *combine-matrix-def*, auto)
apply (rule *comb* [of *Abs-matrix* *Abs-matrix*])
by (auto, insert *combine-infmatrix-assoc*[of *f*], simp add: *associative-def*)

lemma *Rep-apply-matrix*[simp]: $f \ 0 = 0 \implies \text{Rep-matrix } (\text{apply-matrix } f \ A) \ j \ i = f \ (\text{Rep-matrix } A \ j \ i)$
by (simp add: *apply-matrix-def*)

lemma *Rep-combine-matrix*[simp]: $f \ 0 \ 0 = 0 \implies \text{Rep-matrix } (\text{combine-matrix } f \ A \ B) \ j \ i = f \ (\text{Rep-matrix } A \ j \ i) \ (\text{Rep-matrix } B \ j \ i)$
by(simp add: *combine-matrix-def*)

lemma *combine-nrows*: $f \ 0 \ 0 = 0 \implies \text{nrows } (\text{combine-matrix } f \ A \ B) \leq \max (\text{nrows } A) \ (\text{nrows } B)$
by (simp add: *nrows-le*)

lemma *combine-ncols*: $f\ 0\ 0 = 0 \implies \text{ncols}(\text{combine-matrix}\ f\ A\ B) \leq \max(\text{ncols}\ A)\ (\text{ncols}\ B)$

by (*simp add: ncols-le*)

lemma *combine-nrows*: $f\ 0\ 0 = 0 \implies \text{nrows}\ A \leq q \implies \text{nrows}\ B \leq q \implies \text{nrows}(\text{combine-matrix}\ f\ A\ B) \leq q$

by (*simp add: nrows-le*)

lemma *combine-ncols*: $f\ 0\ 0 = 0 \implies \text{ncols}\ A \leq q \implies \text{ncols}\ B \leq q \implies \text{ncols}(\text{combine-matrix}\ f\ A\ B) \leq q$

by (*simp add: ncols-le*)

constdefs

zero-r-neutral :: $('a \Rightarrow 'b::\text{zero} \Rightarrow 'a) \Rightarrow \text{bool}$

zero-r-neutral $f == ! a. f\ a\ 0 = a$

zero-l-neutral :: $('a::\text{zero} \Rightarrow 'b \Rightarrow 'a) \Rightarrow \text{bool}$

zero-l-neutral $f == ! a. f\ 0\ a = a$

zero-closed :: $(('a::\text{zero}) \Rightarrow ('b::\text{zero}) \Rightarrow ('c::\text{zero})) \Rightarrow \text{bool}$

zero-closed $f == (!x. f\ x\ 0 = 0) \ \&\ (!y. f\ 0\ y = 0)$

consts *foldseq* :: $('a \Rightarrow 'a \Rightarrow 'a) \Rightarrow (\text{nat} \Rightarrow 'a) \Rightarrow \text{nat} \Rightarrow 'a$

primrec

foldseq $f\ s\ 0 = s\ 0$

foldseq $f\ s\ (\text{Suc}\ n) = f\ (s\ 0)\ (\text{foldseq}\ f\ (\% k. s(\text{Suc}\ k))\ n)$

consts *foldseq-transposed* :: $('a \Rightarrow 'a \Rightarrow 'a) \Rightarrow (\text{nat} \Rightarrow 'a) \Rightarrow \text{nat} \Rightarrow 'a$

primrec

foldseq-transposed $f\ s\ 0 = s\ 0$

foldseq-transposed $f\ s\ (\text{Suc}\ n) = f\ (\text{foldseq-transposed}\ f\ s\ n)\ (s\ (\text{Suc}\ n))$

lemma *foldseq-assoc* : *associative* $f \implies \text{foldseq}\ f = \text{foldseq-transposed}\ f$

proof –

assume *a:associative* f

then have *sublemma*: $!! n. ! N\ s. N \leq n \longrightarrow \text{foldseq}\ f\ s\ N = \text{foldseq-transposed}\ f\ s\ N$

proof –

fix n

show $!N\ s. N \leq n \longrightarrow \text{foldseq}\ f\ s\ N = \text{foldseq-transposed}\ f\ s\ N$

proof (*induct* n)

show $!N\ s. N \leq 0 \longrightarrow \text{foldseq}\ f\ s\ N = \text{foldseq-transposed}\ f\ s\ N$ **by** *simp*

next

fix n

assume $b: !N\ s. N \leq n \longrightarrow \text{foldseq}\ f\ s\ N = \text{foldseq-transposed}\ f\ s\ N$

have $c: !N\ s. N \leq n \implies \text{foldseq}\ f\ s\ N = \text{foldseq-transposed}\ f\ s\ N$ **by** (*simp add: b*)

show $!N\ t. N \leq \text{Suc}\ n \longrightarrow \text{foldseq}\ f\ t\ N = \text{foldseq-transposed}\ f\ t\ N$

proof (*auto*)

fix $N\ t$

assume $N\text{succ}: N \leq \text{Suc}\ n$

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show foldseq f t N = foldseq-transposed f t N
proof cases
  assume N <= n
  then show foldseq f t N = foldseq-transposed f t N by (simp add: b)
next
  assume ~(N <= n)
  with Nsuc have Nsucq: N = Suc n by simp
  have negz: n ≠ 0 ⇒ ? m. n = Suc m & Suc m <= n by arith
  have assocf: !! x y z. f x (f y z) = f (f x y) z by (insert a, simp add:
associative-def)
  show foldseq f t N = foldseq-transposed f t N
  apply (simp add: Nsucq)
  apply (subst c)
  apply (simp)
  apply (case-tac n = 0)
  apply (simp)
  apply (drule negz)
  apply (erule exE)
  apply (simp)
  apply (subst assocf)
  proof -
    fix m
    assume n = Suc m & Suc m <= n
    then have mless: Suc m <= n by arith
    then have step1: foldseq-transposed f (% k. t (Suc k)) m = foldseq f
(% k. t (Suc k)) m (is ?T1 = ?T2)
      apply (subst c)
      by simp+
    have step2: f (t 0) ?T2 = foldseq f t (Suc m) (is - = ?T3) by simp
    have step3: ?T3 = foldseq-transposed f t (Suc m) (is - = ?T4)
      apply (subst c)
      by (simp add: mless)+
    have step4: ?T4 = f (foldseq-transposed f t m) (t (Suc m)) (is -- ?T5)
by simp
    from step1 step2 step3 step4 show sowhat: f (f (t 0) ?T1) (t (Suc
(Suc m))) = f ?T5 (t (Suc (Suc m))) by simp
    qed
  qed
  qed
  qed
  show foldseq f = foldseq-transposed f by ((rule ext)+, insert sublemma, auto)
qed

lemma foldseq-distr: [[associative f; commutative f]] ⇒ foldseq f (% k. f (u k) (v
k)) n = f (foldseq f u n) (foldseq f v n)
proof -
  assume assoc: associative f
  assume comm: commutative f

```

```

from assoc have a:!!  $x\ y\ z.\ f\ (f\ x\ y)\ z = f\ x\ (f\ y\ z)$  by (simp add: associative-def)
from comm have b:!!  $x\ y.\ f\ x\ y = f\ y\ x$  by (simp add: commutative-def)
from assoc comm have c:!!  $x\ y\ z.\ f\ x\ (f\ y\ z) = f\ y\ (f\ x\ z)$  by (simp add: commutative-def associative-def)
have !! n. (! u v.  $foldseq\ f\ (\%k.\ f\ (u\ k)\ (v\ k))\ n = f\ (foldseq\ f\ u\ n)\ (foldseq\ f\ v\ n)$ )
  apply (induct-tac n)
  apply (simp+, auto)
  by (simp add: a b c)
then show  $foldseq\ f\ (\%k.\ f\ (u\ k)\ (v\ k))\ n = f\ (foldseq\ f\ u\ n)\ (foldseq\ f\ v\ n)$  by
simp
qed

```

```

theorem [associative f; associative g;  $\forall a\ b\ c\ d.\ g\ (f\ a\ b)\ (f\ c\ d) = f\ (g\ a\ c)\ (g\ b\ d)$ ; ? x y.  $(f\ x) \neq (f\ y)$ ; ? x y.  $(g\ x) \neq (g\ y)$ ;  $f\ x\ x = x$ ;  $g\ x\ x = x$ ]  $\implies f=g \mid (!\ y.\ f\ y\ x = y) \mid (!\ y.\ g\ y\ x = y)$ 
oops

```

```

lemma foldseq-zero:
assumes fz:  $f\ 0\ 0 = 0$  and sz: ! i.  $i \leq n \implies s\ i = 0$ 
shows  $foldseq\ f\ s\ n = 0$ 
proof -
  have !! n. ! s. (! i.  $i \leq n \implies s\ i = 0$ )  $\implies foldseq\ f\ s\ n = 0$ 
    apply (induct-tac n)
    apply (simp)
    by (simp add: fz)
  then show  $foldseq\ f\ s\ n = 0$  by (simp add: sz)
qed

```

```

lemma foldseq-significant-positions:
assumes p: ! i.  $i \leq N \implies S\ i = T\ i$ 
shows  $foldseq\ f\ S\ N = foldseq\ f\ T\ N$  (is ?concl)
proof -
  have !! m . ! s t. (! i.  $i \leq m \implies s\ i = t\ i$ )  $\implies foldseq\ f\ s\ m = foldseq\ f\ t\ m$ 
    apply (induct-tac m)
    apply (simp)
    apply (simp)
    apply (auto)
  proof -
    fix n
    fix s::nat $\implies'$ a
    fix t::nat $\implies'$ a
    assume a:  $\forall s\ t.\ (\forall i \leq n.\ s\ i = t\ i) \implies foldseq\ f\ s\ n = foldseq\ f\ t\ n$ 
    assume b:  $\forall i \leq Suc\ n.\ s\ i = t\ i$ 
    have c:!! a b.  $a = b \implies f\ (t\ 0)\ a = f\ (t\ 0)\ b$  by blast
    have d:!! s t. ( $\forall i \leq n.\ s\ i = t\ i$ )  $\implies foldseq\ f\ s\ n = foldseq\ f\ t\ n$  by (simp add: a)
    show  $f\ (t\ 0)\ (foldseq\ f\ (\lambda k.\ s\ (Suc\ k))\ n) = f\ (t\ 0)\ (foldseq\ f\ (\lambda k.\ t\ (Suc$ 

```

k) n) **by** (*rule c, simp add: d b*)
qed
with p **show** $?concl$ **by** *simp*
qed

lemma *foldseq-tail*: $M \leq N \implies \text{foldseq } f \ S \ N = \text{foldseq } f \ (\% \ k. \ (\text{if } k < M \ \text{then } (S \ k) \ \text{else } (\text{foldseq } f \ (\% \ k. \ S(k+M)) \ (N-M)))) \ M$ (**is** $?p \implies ?concl$)

proof –

have *suc*: $!! \ a \ b. \ [a \leq \text{Suc } b; \ a \neq \text{Suc } b] \implies a \leq b$ **by** *arith*
have *a*: $!! \ a \ b \ c. \ a = b \implies f \ c \ a = f \ c \ b$ **by** *blast*
have $!! \ n. \ ! \ m \ s. \ m \leq n \longrightarrow \text{foldseq } f \ s \ n = \text{foldseq } f \ (\% \ k. \ (\text{if } k < m \ \text{then } (s \ k) \ \text{else } (\text{foldseq } f \ (\% \ k. \ s(k+m)) \ (n-m)))) \ m$
apply (*induct-tac n*)
apply (*simp*)
apply (*simp*)
apply (*auto*)
apply (*case-tac m = Suc na*)
apply (*simp*)
apply (*rule a*)
apply (*rule foldseq-significant-positions*)
apply (*auto*)
apply (*drule suc, simp+*)
proof –
fix $na \ m \ s$
assume *suba*: $\forall m \leq na. \ \forall s. \ \text{foldseq } f \ s \ na = \text{foldseq } f \ (\lambda k. \ \text{if } k < m \ \text{then } s \ k \ \text{else } \text{foldseq } f \ (\lambda k. \ s \ (k + m)) \ (na - m)) \ m$
assume *subb*: $m \leq na$
from *suba* **have** *subc*: $!! \ m \ s. \ m \leq na \implies \text{foldseq } f \ s \ na = \text{foldseq } f \ (\lambda k. \ \text{if } k < m \ \text{then } s \ k \ \text{else } \text{foldseq } f \ (\lambda k. \ s \ (k + m)) \ (na - m)) \ m$ **by** *simp*
have *subd*: $\text{foldseq } f \ (\lambda k. \ \text{if } k < m \ \text{then } s \ (\text{Suc } k) \ \text{else } \text{foldseq } f \ (\lambda k. \ s \ (\text{Suc } (k + m))) \ (na - m)) \ m =$
 $\text{foldseq } f \ (\% \ k. \ s \ (\text{Suc } k)) \ na$
by (*rule subc[of m % k. s(Suc k), THEN sym], simp add: subb*)
from *subb* **have** *sube*: $m \neq 0 \implies ? \ mm. \ m = \text{Suc } mm \ \& \ mm \leq na$ **by** *arith*
show $f \ (s \ 0) \ (\text{foldseq } f \ (\lambda k. \ \text{if } k < m \ \text{then } s \ (\text{Suc } k) \ \text{else } \text{foldseq } f \ (\lambda k. \ s \ (\text{Suc } (k + m))) \ (na - m)) \ m) =$
 $\text{foldseq } f \ (\lambda k. \ \text{if } k < m \ \text{then } s \ k \ \text{else } \text{foldseq } f \ (\lambda k. \ s \ (k + m)) \ (\text{Suc } na - m)) \ m$
apply (*simp add: subd*)
apply (*case-tac m=0*)
apply (*simp*)
apply (*drule sube*)
apply (*auto*)
apply (*rule a*)
by (*simp add: subc if-def*)
qed
then **show** $?p \implies ?concl$ **by** *simp*
qed

lemma *foldseq-zerotail*:

assumes

fz: $f\ 0\ 0 = 0$

and *sz*: $! i. n \leq i \longrightarrow s\ i = 0$

and *nm*: $n \leq m$

shows

$foldseq\ f\ s\ n = foldseq\ f\ s\ m$

proof –

show $foldseq\ f\ s\ n = foldseq\ f\ s\ m$

apply (*simp add: foldseq-tail[OF nm, of f s]*)

apply (*rule foldseq-significant-positions*)

apply (*auto*)

apply (*subst foldseq-zero*)

by (*simp add: fz sz*)+

qed

lemma *foldseq-zerotail2*:

assumes $! x. f\ x\ 0 = x$

and $! i. n < i \longrightarrow s\ i = 0$

and *nm*: $n \leq m$

shows

$foldseq\ f\ s\ n = foldseq\ f\ s\ m$ (**is** *?concl*)

proof –

have $f\ 0\ 0 = 0$ **by** (*simp add: prems*)

have $b: !! m\ n. n \leq m \implies m \neq n \implies ? k. m - n = Suc\ k$ **by** *arith*

have $c: 0 \leq m$ **by** *simp*

have $d: !! k. k \neq 0 \implies ? l. k = Suc\ l$ **by** *arith*

show *?concl*

apply (*subst foldseq-tail[OF nm]*)

apply (*rule foldseq-significant-positions*)

apply (*auto*)

apply (*case-tac m=n*)

apply (*simp+*)

apply (*drule b[OF nm]*)

apply (*auto*)

apply (*case-tac k=0*)

apply (*simp add: prems*)

apply (*drule d*)

apply (*auto*)

by (*simp add: prems foldseq-zero*)

qed

lemma *foldseq-zerostart*:

$! x. f\ 0\ (f\ 0\ x) = f\ 0\ x \implies ! i. i \leq n \longrightarrow s\ i = 0 \implies foldseq\ f\ s\ (Suc\ n) = f\ 0\ (s\ (Suc\ n))$

proof –

assume *f00x*: $! x. f\ 0\ (f\ 0\ x) = f\ 0\ x$

have $! s. (! i. i \leq n \longrightarrow s\ i = 0) \longrightarrow foldseq\ f\ s\ (Suc\ n) = f\ 0\ (s\ (Suc\ n))$

```

apply (induct n)
apply (simp)
apply (rule allI, rule impI)
proof -
  fix n
  fix s
  have a:foldseq f s (Suc (Suc n)) = f (s 0) (foldseq f (% k. s(Suc k)) (Suc
n)) by simp
  assume b:!s. (( $\forall i \leq n. s i = 0$ )  $\longrightarrow$  foldseq f s (Suc n) = f 0 (s (Suc n)))
  from b have c:!s. ( $\forall i \leq n. s i = 0$ )  $\implies$  foldseq f s (Suc n) = f 0 (s (Suc
n)) by simp
  assume d:!i. i <= Suc n  $\longrightarrow$  s i = 0
  show foldseq f s (Suc (Suc n)) = f 0 (s (Suc (Suc n)))
    apply (subst a)
    apply (subst c)
    by (simp add: d f00x)+
  qed
  then show !i. i <= n  $\longrightarrow$  s i = 0  $\implies$  foldseq f s (Suc n) = f 0 (s (Suc n))
by simp
qed

```

lemma *foldseq-zerostart2*:

```

!x. f 0 x = x  $\implies$  !i. i < n  $\longrightarrow$  s i = 0  $\implies$  foldseq f s n = s n
proof -
  assume a:!i. i < n  $\longrightarrow$  s i = 0
  assume x:!x. f 0 x = x
  from x have f00x:!x. f 0 (f 0 x) = f 0 x by blast
  have b:!l. i < Suc l = (i <= l) by arith
  have d:!k. k  $\neq$  0  $\implies$  ?l. k = Suc l by arith
  show foldseq f s n = s n
  apply (case-tac n=0)
  apply (simp)
  apply (insert a)
  apply (drule d)
  apply (auto)
  apply (simp add: b)
  apply (insert f00x)
  apply (drule foldseq-zerostart)
  by (simp add: x)+
qed

```

lemma *foldseq-almostzero*:

```

assumes f0x:!x. f 0 x = x and fx0:!x. f x 0 = x and s0:!i. i  $\neq$  j  $\longrightarrow$  s i = 0
shows foldseq f s n = (if (j <= n) then (s j) else 0) (is ?concl)
proof -
  from s0 have a:!i. i < j  $\longrightarrow$  s i = 0 by simp
  from s0 have b:!i. j < i  $\longrightarrow$  s i = 0 by simp
  show ?concl
    apply auto

```

```

apply (subst foldseq-zerotail2[of f, OF fx0, of j, OF b, of n, THEN sym])
apply simp
apply (subst foldseq-zerostart2)
apply (simp add: f0x a)+
apply (subst foldseq-zero)
by (simp add: s0 f0x)+
qed

```

```

lemma foldseq-distr-unary:
  assumes !! a b. g (f a b) = f (g a) (g b)
  shows g(foldseq f s n) = foldseq f (% x. g(s x)) n (is ?concl)
proof -
  have ! s. g(foldseq f s n) = foldseq f (% x. g(s x)) n
    apply (induct-tac n)
    apply (simp)
    apply (simp)
    apply (auto)
    apply (drule-tac x=% k. s (Suc k) in spec)
    by (simp add: prems)
  then show ?concl by simp
qed

```

```

constdefs
  mult-matrix-n :: nat => (('a::zero) => ('b::zero) => ('c::zero)) => ('c => 'c => 'c)
  => 'a matrix => 'b matrix => 'c matrix
  mult-matrix-n n fmul fadd A B == Abs-matrix(% j i. foldseq fadd (% k. fmul
  (Rep-matrix A j k) (Rep-matrix B k i)) n)
  mult-matrix :: (('a::zero) => ('b::zero) => ('c::zero)) => ('c => 'c => 'c) => 'a
  matrix => 'b matrix => 'c matrix
  mult-matrix fmul fadd A B == mult-matrix-n (max (ncols A) (nrows B)) fmul
  fadd A B

```

```

lemma mult-matrix-n:
  assumes prems: ncols A ≤ n (is ?An) nrows B ≤ n (is ?Bn) fadd 0 0 = 0 fmul
  0 0 = 0
  shows c:mult-matrix fmul fadd A B = mult-matrix-n n fmul fadd A B (is ?concl)
proof -
  show ?concl using prems
    apply (simp add: mult-matrix-def mult-matrix-n-def)
    apply (rule comb[of Abs-matrix Abs-matrix], simp, (rule ext)+)
    by (rule foldseq-zerotail, simp-all add: nrows-le ncols-le prems)
qed

```

```

lemma mult-matrix-nm:
  assumes prems: ncols A ≤ n nrows B ≤ n ncols A ≤ m nrows B ≤ m
  fadd 0 0 = 0 fmul 0 0 = 0
  shows mult-matrix-n n fmul fadd A B = mult-matrix-n m fmul fadd A B
proof -
  from prems have mult-matrix-n n fmul fadd A B = mult-matrix fmul fadd A B

```

by (simp add: mult-matrix-n)
 also from prems have ... = mult-matrix-n m fmul fadd A B by (simp add:
 mult-matrix-n[THEN sym])
 finally show mult-matrix-n n fmul fadd A B = mult-matrix-n m fmul fadd A B
 by simp
 qed

constdefs

r-distributive :: ('a ⇒ 'b ⇒ 'b) ⇒ ('b ⇒ 'b ⇒ 'b) ⇒ bool
 r-distributive fmul fadd == ! a u v. fmul a (fadd u v) = fadd (fmul a u) (fmul a
 v)
 l-distributive :: ('a ⇒ 'b ⇒ 'a) ⇒ ('a ⇒ 'a ⇒ 'a) ⇒ bool
 l-distributive fmul fadd == ! a u v. fmul (fadd u v) a = fadd (fmul u a) (fmul v
 a)
 distributive :: ('a ⇒ 'a ⇒ 'a) ⇒ ('a ⇒ 'a ⇒ 'a) ⇒ bool
 distributive fmul fadd == l-distributive fmul fadd & r-distributive fmul fadd

lemma max1: !! a x y. (a::nat) <= x ⇒ a <= max x y by (arith)

lemma max2: !! b x y. (b::nat) <= y ⇒ b <= max x y by (arith)

lemma r-distributive-matrix:

assumes prems:

r-distributive fmul fadd

associative fadd

commutative fadd

fadd 0 0 = 0

! a. fmul a 0 = 0

! a. fmul 0 a = 0

shows r-distributive (mult-matrix fmul fadd) (combine-matrix fadd) (is ?concl)

proof -

from prems show ?concl

apply (simp add: r-distributive-def mult-matrix-def, auto)

proof -

fix a::'a matrix

fix u::'b matrix

fix v::'b matrix

let ?mx = max (ncols a) (max (nrows u) (nrows v))

from prems show mult-matrix-n (max (ncols a) (nrows (combine-matrix fadd
 u v))) fmul fadd a (combine-matrix fadd u v) =

combine-matrix fadd (mult-matrix-n (max (ncols a) (nrows u)) fmul fadd a
 u) (mult-matrix-n (max (ncols a) (nrows v)) fmul fadd a v)

apply (subst mult-matrix-nm[of - - - ?mx fadd fmul])

apply (simp add: max1 max2 combine-nrows combine-ncols)+

apply (subst mult-matrix-nm[of - - v ?mx fadd fmul])

apply (simp add: max1 max2 combine-nrows combine-ncols)+

apply (subst mult-matrix-nm[of - - u ?mx fadd fmul])

apply (simp add: max1 max2 combine-nrows combine-ncols)+

apply (simp add: mult-matrix-n-def r-distributive-def foldseq-distr[of fadd])

apply (simp add: combine-matrix-def combine-infmatrix-def)

```

apply (rule comb[of Abs-matrix Abs-matrix], simp, (rule ext)+)
apply (simplesubst RepAbs-matrix)
apply (simp, auto)
apply (rule exI[of - nrows a], simp add: nrows-le foldseq-zero)
apply (rule exI[of - ncols v], simp add: ncols-le foldseq-zero)
apply (subst RepAbs-matrix)
apply (simp, auto)
apply (rule exI[of - nrows a], simp add: nrows-le foldseq-zero)
apply (rule exI[of - ncols u], simp add: ncols-le foldseq-zero)
done
qed
qed

lemma l-distributive-matrix:
  assumes prems:
    l-distributive fmul fadd
    associative fadd
    commutative fadd
    fadd 0 0 = 0
    ! a. fmul a 0 = 0
    ! a. fmul 0 a = 0
  shows l-distributive (mult-matrix fmul fadd) (combine-matrix fadd) (is ?concl)
proof -
  from prems show ?concl
  apply (simp add: l-distributive-def mult-matrix-def, auto)
proof -
  fix a::'b matrix
  fix u::'a matrix
  fix v::'a matrix
  let ?mx = max (nrows a) (max (ncols u) (ncols v))
  from prems show mult-matrix-n (max (ncols (combine-matrix fadd u v))
(nrows a)) fmul fadd (combine-matrix fadd u v) a =
    combine-matrix fadd (mult-matrix-n (max (ncols u) (nrows a)) fmul
fadd u a) (mult-matrix-n (max (ncols v) (nrows a)) fmul fadd v a)
  apply (subst mult-matrix-nm[of v - - ?mx fadd fmul])
  apply (simp add: max1 max2 combine-nrows combine-ncols)+
  apply (subst mult-matrix-nm[of u - - ?mx fadd fmul])
  apply (simp add: max1 max2 combine-nrows combine-ncols)+
  apply (subst mult-matrix-nm[of - - - ?mx fadd fmul])
  apply (simp add: max1 max2 combine-nrows combine-ncols)+
  apply (simp add: mult-matrix-n-def l-distributive-def foldseq-distr[of fadd])
  apply (simp add: combine-matrix-def combine-infmatrix-def)
  apply (rule comb[of Abs-matrix Abs-matrix], simp, (rule ext)+)
  apply (simplesubst RepAbs-matrix)
  apply (simp, auto)
  apply (rule exI[of - nrows v], simp add: nrows-le foldseq-zero)
  apply (rule exI[of - ncols a], simp add: ncols-le foldseq-zero)
  apply (subst RepAbs-matrix)
  apply (simp, auto)

```

```

    apply (rule exI[of - nrows u], simp add: nrows-le foldseq-zero)
    apply (rule exI[of - ncols a], simp add: ncols-le foldseq-zero)
  done
qed
qed

instance matrix :: (zero) zero ..

defs(overloaded)
  zero-matrix-def: (0::('a::zero) matrix) == Abs-matrix(% j i. 0)

lemma Rep-zero-matrix-def[simp]: Rep-matrix 0 j i = 0
  apply (simp add: zero-matrix-def)
  apply (subst RepAbs-matrix)
  by (auto)

lemma zero-matrix-def-nrows[simp]: nrows 0 = 0
proof -
  have a!! (x::nat). x <= 0 ==> x = 0 by (arith)
  show nrows 0 = 0 by (rule a, subst nrows-le, simp)
qed

lemma zero-matrix-def-ncols[simp]: ncols 0 = 0
proof -
  have a!! (x::nat). x <= 0 ==> x = 0 by (arith)
  show ncols 0 = 0 by (rule a, subst ncols-le, simp)
qed

lemma combine-matrix-zero-l-neutral: zero-l-neutral f ==> zero-l-neutral (combine-matrix
f)
  by (simp add: zero-l-neutral-def combine-matrix-def combine-infmatrix-def)

lemma combine-matrix-zero-r-neutral: zero-r-neutral f ==> zero-r-neutral (combine-matrix
f)
  by (simp add: zero-r-neutral-def combine-matrix-def combine-infmatrix-def)

lemma mult-matrix-zero-closed: [[fadd 0 0 = 0; zero-closed fmul]] ==> zero-closed
(mult-matrix fmul fadd)
  apply (simp add: zero-closed-def mult-matrix-def mult-matrix-n-def)
  apply (auto)
  by (subst foldseq-zero, (simp add: zero-matrix-def)+)+

lemma mult-matrix-n-zero-right[simp]: [[fadd 0 0 = 0; !a. fmul a 0 = 0]] ==>
mult-matrix-n n fmul fadd A 0 = 0
  apply (simp add: mult-matrix-n-def)
  apply (subst foldseq-zero)
  by (simp-all add: zero-matrix-def)

lemma mult-matrix-n-zero-left[simp]: [[fadd 0 0 = 0; !a. fmul 0 a = 0]] ==>

```

```

mult-matrix-n n fmul fadd 0 A = 0
apply (simp add: mult-matrix-n-def)
apply (subst foldseq-zero)
by (simp-all add: zero-matrix-def)

```

```

lemma mult-matrix-zero-left[simp]:  $\llbracket fadd\ 0\ 0 = 0; !a.\ fmul\ 0\ a = 0 \rrbracket \implies$  mult-matrix
fmul fadd 0 A = 0
by (simp add: mult-matrix-def)

```

```

lemma mult-matrix-zero-right[simp]:  $\llbracket fadd\ 0\ 0 = 0; !a.\ fmul\ a\ 0 = 0 \rrbracket \implies$  mult-matrix
fmul fadd A 0 = 0
by (simp add: mult-matrix-def)

```

```

lemma apply-matrix-zero[simp]:  $f\ 0 = 0 \implies$  apply-matrix  $f\ 0 = 0$ 
apply (simp add: apply-matrix-def apply-infmatrix-def)
by (simp add: zero-matrix-def)

```

```

lemma combine-matrix-zero:  $f\ 0\ 0 = 0 \implies$  combine-matrix  $f\ 0\ 0 = 0$ 
apply (simp add: combine-matrix-def combine-infmatrix-def)
by (simp add: zero-matrix-def)

```

```

lemma transpose-matrix-zero[simp]: transpose-matrix  $0 = 0$ 
apply (simp add: transpose-matrix-def transpose-infmatrix-def zero-matrix-def RepAbs-matrix)
apply (subst Rep-matrix-inject[symmetric], (rule ext)+)
apply (simp add: RepAbs-matrix)
done

```

```

lemma apply-zero-matrix-def[simp]: apply-matrix  $(\% x.\ 0)\ A = 0$ 
apply (simp add: apply-matrix-def apply-infmatrix-def)
by (simp add: zero-matrix-def)

```

constdefs

```

singleton-matrix ::  $nat \Rightarrow nat \Rightarrow ('a::zero) \Rightarrow 'a\ matrix$ 
singleton-matrix  $j\ i\ a == Abs-matrix(\% m\ n.\ if\ j = m \ \&\ i = n\ then\ a\ else\ 0)$ 
move-matrix ::  $('a::zero)\ matrix \Rightarrow int \Rightarrow int \Rightarrow 'a\ matrix$ 
move-matrix  $A\ y\ x == Abs-matrix(\% j\ i.\ if\ (neg\ ((int\ j)-y)) \ | (neg\ ((int\ i)-x))$ 
then  $0\ else\ Rep-matrix\ A\ (nat\ ((int\ j)-y))\ (nat\ ((int\ i)-x))$ )
take-rows ::  $('a::zero)\ matrix \Rightarrow nat \Rightarrow 'a\ matrix$ 
take-rows  $A\ r == Abs-matrix(\% j\ i.\ if\ (j < r)\ then\ (Rep-matrix\ A\ j\ i)\ else\ 0)$ 
take-columns ::  $('a::zero)\ matrix \Rightarrow nat \Rightarrow 'a\ matrix$ 
take-columns  $A\ c == Abs-matrix(\% j\ i.\ if\ (i < c)\ then\ (Rep-matrix\ A\ j\ i)\ else$ 
 $0)$ 

```

constdefs

```

column-of-matrix ::  $('a::zero)\ matrix \Rightarrow nat \Rightarrow 'a\ matrix$ 
column-of-matrix  $A\ n == take-columns\ (move-matrix\ A\ 0\ (-\ int\ n))\ 1$ 
row-of-matrix ::  $('a::zero)\ matrix \Rightarrow nat \Rightarrow 'a\ matrix$ 
row-of-matrix  $A\ m == take-rows\ (move-matrix\ A\ (-\ int\ m)\ 0)\ 1$ 

```

lemma *Rep-singleton-matrix*[simp]: *Rep-matrix* (*singleton-matrix* *j i e*) *m n* = (if *j = m* & *i = n* then *e* else 0)
apply (*simp add: singleton-matrix-def*)
apply (*auto*)
apply (*subst RepAbs-matrix*)
apply (*rule exI[of - Suc m], simp*)
apply (*rule exI[of - Suc n], simp+*)
by (*subst RepAbs-matrix, rule exI[of - Suc j], simp, rule exI[of - Suc i], simp+*)+

lemma *apply-singleton-matrix*[simp]: $f\ 0 = 0 \implies \text{apply-matrix } f \text{ (singleton-matrix } j\ i\ x) = \text{singleton-matrix } j\ i\ (f\ x)$
apply (*subst Rep-matrix-inject[symmetric]*)
apply (*rule ext*)
apply (*simp*)
done

lemma *singleton-matrix-zero*[simp]: *singleton-matrix* *j i 0* = 0
by (*simp add: singleton-matrix-def zero-matrix-def*)

lemma *nrows-singleton*[simp]: $\text{nrows}(\text{singleton-matrix } j\ i\ e) = (\text{if } e = 0 \text{ then } 0 \text{ else } \text{Suc } j)$
apply (*auto*)
apply (*rule le-anti-sym*)
apply (*subst nrows-le*)
apply (*simp add: singleton-matrix-def, auto*)
apply (*subst RepAbs-matrix*)
apply (*simp, arith*)
apply (*simp, arith*)
apply (*simp*)
apply (*simp add: Suc-le-eq*)
apply (*rule not-leE*)
apply (*subst nrows-le*)
by *simp*

lemma *ncols-singleton*[simp]: $\text{ncols}(\text{singleton-matrix } j\ i\ e) = (\text{if } e = 0 \text{ then } 0 \text{ else } \text{Suc } i)$
apply (*auto*)
apply (*rule le-anti-sym*)
apply (*subst ncols-le*)
apply (*simp add: singleton-matrix-def, auto*)
apply (*subst RepAbs-matrix*)
apply (*simp, arith*)
apply (*simp, arith*)
apply (*simp*)
apply (*simp add: Suc-le-eq*)
apply (*rule not-leE*)
apply (*subst ncols-le*)
by *simp*

lemma *combine-singleton*: $f\ 0\ 0 = 0 \implies \text{combine-matrix } f\ (\text{singleton-matrix } j\ i\ a)\ (\text{singleton-matrix } j\ i\ b) = \text{singleton-matrix } j\ i\ (f\ a\ b)$
apply (*simp add: singleton-matrix-def combine-matrix-def combine-infmatrix-def*)
apply (*subst RepAbs-matrix*)
apply (*rule exI[of - Suc j], simp*)
apply (*rule exI[of - Suc i], simp*)
apply (*rule comb[of Abs-matrix Abs-matrix], simp, (rule ext)+*)
apply (*subst RepAbs-matrix*)
apply (*rule exI[of - Suc j], simp*)
apply (*rule exI[of - Suc i], simp*)
by *simp*

lemma *transpose-singleton[simp]*: $\text{transpose-matrix } (\text{singleton-matrix } j\ i\ a) = \text{singleton-matrix } i\ j\ a$
apply (*subst Rep-matrix-inject[symmetric], (rule ext)+*)
apply (*simp*)
done

lemma *Rep-move-matrix[simp]*:
 $\text{Rep-matrix } (\text{move-matrix } A\ y\ x)\ j\ i =$
 $(\text{if } (\text{neg } ((\text{int } j) - y)) \mid (\text{neg } ((\text{int } i) - x)) \text{ then } 0 \text{ else } \text{Rep-matrix } A\ (\text{nat}((\text{int } j) - y))$
 $(\text{nat}((\text{int } i) - x)))$
apply (*simp add: move-matrix-def*)
apply (*auto*)
by (*subst RepAbs-matrix,*
rule exI[of - (nrows A)+(nat (abs y))], auto, rule nrows, arith,
rule exI[of - (ncols A)+(nat (abs x))], auto, rule ncols, arith)+

lemma *move-matrix-0-0[simp]*: $\text{move-matrix } A\ 0\ 0 = A$
by (*simp add: move-matrix-def*)

lemma *move-matrix-ortho*: $\text{move-matrix } A\ j\ i = \text{move-matrix } (\text{move-matrix } A\ j\ 0)\ 0\ i$
apply (*subst Rep-matrix-inject[symmetric]*)
apply (*rule ext*)
apply (*simp*)
done

lemma *transpose-move-matrix[simp]*:
 $\text{transpose-matrix } (\text{move-matrix } A\ x\ y) = \text{move-matrix } (\text{transpose-matrix } A)\ y\ x$
apply (*subst Rep-matrix-inject[symmetric], (rule ext)+*)
apply (*simp*)
done

lemma *move-matrix-singleton[simp]*: $\text{move-matrix } (\text{singleton-matrix } u\ v\ x)\ j\ i =$
 $(\text{if } (j + \text{int } u < 0) \mid (i + \text{int } v < 0) \text{ then } 0 \text{ else } (\text{singleton-matrix } (\text{nat } (j + \text{int } u))\ (\text{nat } (i + \text{int } v))\ x))$
apply (*subst Rep-matrix-inject[symmetric]*)
apply (*rule ext*)
done

```

apply (case-tac  $j + \text{int } u < 0$ )
apply (simp, arith)
apply (case-tac  $i + \text{int } v < 0$ )
apply (simp add: neg-def, arith)
apply (simp add: neg-def)
apply arith
done

```

```

lemma Rep-take-columns[simp]:
  Rep-matrix (take-columns  $A$   $c$ )  $j$   $i$  =
    (if  $i < c$  then (Rep-matrix  $A$   $j$   $i$ ) else 0)
apply (simp add: take-columns-def)
apply (simplesubst RepAbs-matrix)
apply (rule exI[of - nrows  $A$ ], auto, simp add: nrows-le)
apply (rule exI[of - ncols  $A$ ], auto, simp add: ncols-le)
done

```

```

lemma Rep-take-rows[simp]:
  Rep-matrix (take-rows  $A$   $r$ )  $j$   $i$  =
    (if  $j < r$  then (Rep-matrix  $A$   $j$   $i$ ) else 0)
apply (simp add: take-rows-def)
apply (simplesubst RepAbs-matrix)
apply (rule exI[of - nrows  $A$ ], auto, simp add: nrows-le)
apply (rule exI[of - ncols  $A$ ], auto, simp add: ncols-le)
done

```

```

lemma Rep-column-of-matrix[simp]:
  Rep-matrix (column-of-matrix  $A$   $c$ )  $j$   $i$  = (if  $i = 0$  then (Rep-matrix  $A$   $j$   $c$ ) else
0)
  by (simp add: column-of-matrix-def)

```

```

lemma Rep-row-of-matrix[simp]:
  Rep-matrix (row-of-matrix  $A$   $r$ )  $j$   $i$  = (if  $j = 0$  then (Rep-matrix  $A$   $r$   $i$ ) else 0)
  by (simp add: row-of-matrix-def)

```

```

lemma column-of-matrix: ncols  $A$   $\leq n \implies$  column-of-matrix  $A$   $n = 0$ 
apply (subst Rep-matrix-inject[THEN sym])
apply (rule ext)+
by (simp add: ncols)

```

```

lemma row-of-matrix: nrows  $A$   $\leq n \implies$  row-of-matrix  $A$   $n = 0$ 
apply (subst Rep-matrix-inject[THEN sym])
apply (rule ext)+
by (simp add: nrows)

```

```

lemma mult-matrix-singleton-right[simp]:
  assumes prems:
    !  $x$ . fmul  $x$  0 = 0
    !  $x$ . fmul 0  $x$  = 0

```

```

! x. fadd 0 x = x
! x. fadd x 0 = x
shows (mult-matrix fmul fadd A (singleton-matrix j i e)) = apply-matrix (% x.
fmul x e) (move-matrix (column-of-matrix A j) 0 (int i))
apply (simp add: mult-matrix-def)
apply (subst mult-matrix-nm[of - - - max (ncols A) (Suc j)])
apply (simp add: max-def)+
apply (auto)
apply (simp add: prems)+
apply (simp add: mult-matrix-n-def apply-matrix-def apply-infmatrix-def)
apply (rule comb[of Abs-matrix Abs-matrix], auto, (rule ext)+)
apply (subst foldseq-almostzero[of - j])
apply (simp add: prems)+
apply (auto)
apply (insert ncols-le[of A j])
apply (arith)
proof -
  fix k
  fix l
  assume a: ~neg(int l - int i)
  assume b: nat (int l - int i) = 0
  from a b have a: l = i by (insert not-neg-nat[of int l - int i], simp add: a b)
  assume c: i ≠ l
  from c a show fmul (Rep-matrix A k j) e = 0 by blast
qed

```

lemma *mult-matrix-ext*:

```

assumes
  eprem:
  ? e. (! a b. a ≠ b → fmul a e ≠ fmul b e)
and fpreds:
  ! a. fmul 0 a = 0
  ! a. fmul a 0 = 0
  ! a. fadd a 0 = a
  ! a. fadd 0 a = a
and contrapreds:
  mult-matrix fmul fadd A = mult-matrix fmul fadd B
shows
  A = B
proof (rule contrapos-np[of False], simp)
  assume a: A ≠ B
  have b: !! f g. (! x y. f x y = g x y) ⇒ f = g by ((rule ext)+, auto)
  have ? j i. (Rep-matrix A j i) ≠ (Rep-matrix B j i)
    apply (rule contrapos-np[of False], simp+)
    apply (insert b[of Rep-matrix A Rep-matrix B], simp)
    by (simp add: Rep-matrix-inject a)
  then obtain J I where c: (Rep-matrix A J I) ≠ (Rep-matrix B J I) by blast
  from eprem obtain e where eprops: (! a b. a ≠ b → fmul a e ≠ fmul b e) by
  blast

```

```

let ?S = singleton-matrix I 0 e
let ?comp = mult-matrix fmul fadd
have d: !!x f g. f = g  $\implies$  f x = g x by blast
have e: (% x. fmul x e) 0 = 0 by (simp add: prems)
have ~(?comp A ?S = ?comp B ?S)
  apply (rule notI)
  apply (simp add: fprems eprops)
  apply (simp add: Rep-matrix-inject[THEN sym])
  apply (drule d[of - - J], drule d[of - - 0])
  by (simp add: e c eprops)
with contrapremis show False by simp
qed

```

constdefs

```

foldmatrix :: ('a  $\Rightarrow$  'a  $\Rightarrow$  'a)  $\Rightarrow$  ('a  $\Rightarrow$  'a  $\Rightarrow$  'a)  $\Rightarrow$  ('a infmatrix)  $\Rightarrow$  nat  $\Rightarrow$  nat
 $\Rightarrow$  'a
foldmatrix f g A m n == foldseq-transposed g (% j. foldseq f (A j) n) m
foldmatrix-transposed :: ('a  $\Rightarrow$  'a  $\Rightarrow$  'a)  $\Rightarrow$  ('a  $\Rightarrow$  'a  $\Rightarrow$  'a)  $\Rightarrow$  ('a infmatrix)  $\Rightarrow$ 
nat  $\Rightarrow$  nat  $\Rightarrow$  'a
foldmatrix-transposed f g A m n == foldseq g (% j. foldseq-transposed f (A j) n)
m

```

lemma foldmatrix-transpose:

```

assumes
! a b c d. g(f a b) (f c d) = f (g a c) (g b d)
shows
foldmatrix f g A m n = foldmatrix-transposed g f (transpose-infmatrix A) n m
(is ?concl)
proof -
have forall:!! P x. (! x. P x)  $\implies$  P x by auto
have tworows:! A. foldmatrix f g A 1 n = foldmatrix-transposed g f (transpose-infmatrix
A) n 1
apply (induct n)
apply (simp add: foldmatrix-def foldmatrix-transposed-def prems)+
apply (auto)
by (drule-tac x=(% j i. A j (Suc i)) in forall, simp)
show foldmatrix f g A m n = foldmatrix-transposed g f (transpose-infmatrix A)
n m
apply (simp add: foldmatrix-def foldmatrix-transposed-def)
apply (induct m, simp)
apply (simp)
apply (insert tworows)
apply (drule-tac x=% j i. (if j = 0 then (foldseq-transposed g ( $\lambda$ u. A u i) m)
else (A (Suc m) i)) in spec)
by (simp add: foldmatrix-def foldmatrix-transposed-def)
qed

```

lemma foldseq-foldseq:

assumes

associative f
associative g
 ! a b c d. $g(f a b) (f c d) = f (g a c) (g b d)$
shows
 foldseq g (% j. foldseq f (A j) n) m = foldseq f (% j. foldseq g ((transpose-infmatrix A) j) m) n
apply (insert foldmatrix-transpose[of g f A m n])
by (simp add: foldmatrix-def foldmatrix-transposed-def foldseq-assoc[THEN sym] prems)

lemma *mult-n-nrows:*

assumes
 ! a. $fmul\ 0\ a = 0$
 ! a. $fmul\ a\ 0 = 0$
 $fadd\ 0\ 0 = 0$
shows $nrows\ (mult-matrix-n\ n\ fmul\ fadd\ A\ B) \leq nrows\ A$
apply (subst nrows-le)
apply (simp add: mult-matrix-n-def)
apply (subst RepAbs-matrix)
apply (rule-tac x=nrows A in exI)
apply (simp add: nrows prems foldseq-zero)
apply (rule-tac x=ncols B in exI)
apply (simp add: ncols prems foldseq-zero)
by (simp add: nrows prems foldseq-zero)

lemma *mult-n-ncols:*

assumes
 ! a. $fmul\ 0\ a = 0$
 ! a. $fmul\ a\ 0 = 0$
 $fadd\ 0\ 0 = 0$
shows $ncols\ (mult-matrix-n\ n\ fmul\ fadd\ A\ B) \leq ncols\ B$
apply (subst ncols-le)
apply (simp add: mult-matrix-n-def)
apply (subst RepAbs-matrix)
apply (rule-tac x=nrows A in exI)
apply (simp add: nrows prems foldseq-zero)
apply (rule-tac x=ncols B in exI)
apply (simp add: ncols prems foldseq-zero)
by (simp add: ncols prems foldseq-zero)

lemma *mult-nrows:*

assumes
 ! a. $fmul\ 0\ a = 0$
 ! a. $fmul\ a\ 0 = 0$
 $fadd\ 0\ 0 = 0$
shows $nrows\ (mult-matrix\ fmul\ fadd\ A\ B) \leq nrows\ A$
by (simp add: mult-matrix-def mult-n-nrows prems)

lemma *mult-ncols:*

```

assumes
! a. fmul 0 a = 0
! a. fmul a 0 = 0
fadd 0 0 = 0
shows ncols (mult-matrix fmul fadd A B) ≤ ncols B
by (simp add: mult-matrix-def mult-n-ncols prems)

lemma nrows-move-matrix-le: nrows (move-matrix A j i) <= nat((int (nrows A))
+ j)
apply (auto simp add: nrows-le)
apply (rule nrows)
apply (arith)
done

lemma ncols-move-matrix-le: ncols (move-matrix A j i) <= nat((int (ncols A))
+ i)
apply (auto simp add: ncols-le)
apply (rule ncols)
apply (arith)
done

lemma mult-matrix-assoc:
assumes prems:
! a. fmul1 0 a = 0
! a. fmul1 a 0 = 0
! a. fmul2 0 a = 0
! a. fmul2 a 0 = 0
fadd1 0 0 = 0
fadd2 0 0 = 0
! a b c d. fadd2 (fadd1 a b) (fadd1 c d) = fadd1 (fadd2 a c) (fadd2 b d)
associative fadd1
associative fadd2
! a b c. fmul2 (fmul1 a b) c = fmul1 a (fmul2 b c)
! a b c. fmul2 (fadd1 a b) c = fadd1 (fmul2 a c) (fmul2 b c)
! a b c. fmul1 c (fadd2 a b) = fadd2 (fmul1 c a) (fmul1 c b)
shows mult-matrix fmul2 fadd2 (mult-matrix fmul1 fadd1 A B) C = mult-matrix
fmul1 fadd1 A (mult-matrix fmul2 fadd2 B C) (is ?concl)
proof –
have comb-left: !! A B x y. A = B ⇒ (Rep-matrix (Abs-matrix A)) x y =
(Rep-matrix(Abs-matrix B)) x y by blast
have fmul2fadd1fold: !! x s n. fmul2 (foldseq fadd1 s n) x = foldseq fadd1 (%
k. fmul2 (s k) x) n
by (rule-tac g1 = % y. fmul2 y x in ssubst [OF foldseq-distr-unary], simp-all!)
have fmul1fadd2fold: !! x s n. fmul1 x (foldseq fadd2 s n) = foldseq fadd2 (% k.
fmul1 x (s k)) n
by (rule-tac g1 = % y. fmul1 x y in ssubst [OF foldseq-distr-unary], simp-all!)
let ?N = max (ncols A) (max (ncols B) (max (nrows B) (nrows C)))
show ?concl
apply (simp add: Rep-matrix-inject[THEN sym])

```

```

apply (rule ext)+
apply (simp add: mult-matrix-def)
apply (simplesubst mult-matrix-nm[of - max (ncols (mult-matrix-n (max (ncols
A) (nrows B)) fmul1 fadd1 A B)) (nrows C) - max (ncols B) (nrows C)])
apply (simp add: max1 max2 mult-n-ncols mult-n-nrows prems)+
apply (simplesubst mult-matrix-nm[of - max (ncols A) (nrows (mult-matrix-n
(max (ncols B) (nrows C)) fmul2 fadd2 B C)) - max (ncols A) (nrows B)]) ap-
ply (simp add: max1 max2 mult-n-ncols mult-n-nrows prems)+
apply (simplesubst mult-matrix-nm[of - - - ?N])
apply (simp add: max1 max2 mult-n-ncols mult-n-nrows prems)+
apply (simplesubst mult-matrix-nm[of - - - ?N])
apply (simp add: max1 max2 mult-n-ncols mult-n-nrows prems)+
apply (simplesubst mult-matrix-nm[of - - - ?N])
apply (simp add: max1 max2 mult-n-ncols mult-n-nrows prems)+
apply (simp add: mult-matrix-n-def)
apply (rule comb-left)
apply ((rule ext)+, simp)
apply (simplesubst RepAbs-matrix)
apply (rule exI[of - nrows B])
apply (simp add: nrows prems foldseq-zero)
apply (rule exI[of - ncols C])
apply (simp add: prems ncols foldseq-zero)
apply (subst RepAbs-matrix)
apply (rule exI[of - nrows A])
apply (simp add: nrows prems foldseq-zero)
apply (rule exI[of - ncols B])
apply (simp add: prems ncols foldseq-zero)
apply (simp add: fmul2fadd1fold fmul1fadd2fold prems)
apply (subst foldseq-foldseq)
apply (simp add: prems)+
by (simp add: transpose-infmatrix)
qed

```

lemma

```

assumes prems:
! a. fmul1 0 a = 0
! a. fmul1 a 0 = 0
! a. fmul2 0 a = 0
! a. fmul2 a 0 = 0
fadd1 0 0 = 0
fadd2 0 0 = 0
! a b c d. fadd2 (fadd1 a b) (fadd1 c d) = fadd1 (fadd2 a c) (fadd2 b d)
associative fadd1
associative fadd2
! a b c. fmul2 (fmul1 a b) c = fmul1 a (fmul2 b c)
! a b c. fmul2 (fadd1 a b) c = fadd1 (fmul2 a c) (fmul2 b c)
! a b c. fmul1 c (fadd2 a b) = fadd2 (fmul1 c a) (fmul1 c b)

```

shows
 $(\text{mult-matrix } \text{fmul1 } \text{fadd1 } A) \circ (\text{mult-matrix } \text{fmul2 } \text{fadd2 } B) = \text{mult-matrix } \text{fmul2 } \text{fadd2 } (\text{mult-matrix } \text{fmul1 } \text{fadd1 } A B)$
apply (rule ext)+
apply (simp add: comp-def)
by (simp add: mult-matrix-assoc prems)

lemma *mult-matrix-assoc-simple*:

assumes *prems*:
! a. $\text{fmul } 0 \ a = 0$
! a. $\text{fmul } a \ 0 = 0$
 $\text{fadd } 0 \ 0 = 0$
associative fadd
commutative fadd
associative fmul
distributive fmul fadd
shows $\text{mult-matrix } \text{fmul } \text{fadd } (\text{mult-matrix } \text{fmul } \text{fadd } A B) C = \text{mult-matrix } \text{fmul } \text{fadd } A (\text{mult-matrix } \text{fmul } \text{fadd } B C)$ (**is** ?concl)
proof –
have !! a b c d. $\text{fadd } (\text{fadd } a \ b) (\text{fadd } c \ d) = \text{fadd } (\text{fadd } a \ c) (\text{fadd } b \ d)$
by (simp! add: associative-def commutative-def)
then show ?concl
apply (subst mult-matrix-assoc)
apply (simp-all!)
by (simp add: associative-def distributive-def l-distributive-def r-distributive-def)+
qed

lemma *transpose-apply-matrix*: $f \ 0 = 0 \implies \text{transpose-matrix } (\text{apply-matrix } f \ A) = \text{apply-matrix } f \ (\text{transpose-matrix } A)$
apply (simp add: Rep-matrix-inject[THEN sym])
apply (rule ext)+
by simp

lemma *transpose-combine-matrix*: $f \ 0 \ 0 = 0 \implies \text{transpose-matrix } (\text{combine-matrix } f \ A \ B) = \text{combine-matrix } f \ (\text{transpose-matrix } A) \ (\text{transpose-matrix } B)$
apply (simp add: Rep-matrix-inject[THEN sym])
apply (rule ext)+
by simp

lemma *Rep-mult-matrix*:

assumes
! a. $\text{fmul } 0 \ a = 0$
! a. $\text{fmul } a \ 0 = 0$
 $\text{fadd } 0 \ 0 = 0$
shows
 $\text{Rep-matrix}(\text{mult-matrix } \text{fmul } \text{fadd } A \ B) \ j \ i = \text{foldseq } \text{fadd } (\% k. \text{fmul } (\text{Rep-matrix } A \ j \ k) (\text{Rep-matrix } B \ k \ i)) (\text{max } (\text{ncols } A) (\text{nrows } B))$
apply (simp add: mult-matrix-def mult-matrix-n-def)

```

apply (subst RepAbs-matrix)
apply (rule exI[of - nrows A], simp! add: nrows foldseq-zero)
apply (rule exI[of - ncols B], simp! add: ncols foldseq-zero)
by simp

```

lemma transpose-mult-matrix:

```

assumes
  ! a. fmul 0 a = 0
  ! a. fmul a 0 = 0
  fadd 0 0 = 0
  ! x y. fmul y x = fmul x y
shows
  transpose-matrix (mult-matrix fmul fadd A B) = mult-matrix fmul fadd (transpose-matrix
  B) (transpose-matrix A)
apply (simp add: Rep-matrix-inject[THEN sym])
apply (rule ext)+
by (simp! add: Rep-mult-matrix max-ac)

```

lemma column-transpose-matrix: column-of-matrix (transpose-matrix A) n = transpose-matrix (row-of-matrix A n)

```

apply (simp add: Rep-matrix-inject[THEN sym])
apply (rule ext)+
by simp

```

lemma take-columns-transpose-matrix: take-columns (transpose-matrix A) n = transpose-matrix (take-rows A n)

```

apply (simp add: Rep-matrix-inject[THEN sym])
apply (rule ext)+
by simp

```

instance matrix :: ({ord, zero}) ord ..

defs (overloaded)

```

  le-matrix-def: (A::('a::{ord,zero}) matrix) <= B == ! j i. (Rep-matrix A j i)
  <= (Rep-matrix B j i)
  less-def: (A::('a::{ord,zero}) matrix) < B == (A <= B) & (A ≠ B)

```

instance matrix :: ({order, zero}) order

```

apply intro-classes
apply (simp-all add: le-matrix-def less-def)
apply (auto)
apply (drule-tac x=j in spec, drule-tac x=j in spec)
apply (drule-tac x=i in spec, drule-tac x=i in spec)
apply (simp)
apply (simp add: Rep-matrix-inject[THEN sym])
apply (rule ext)+
apply (drule-tac x=xa in spec, drule-tac x=xa in spec)
apply (drule-tac x=xb in spec, drule-tac x=xb in spec)
by simp

```

lemma *le-apply-matrix*:

assumes

$f\ 0 = 0$

$! x\ y. x \leq y \longrightarrow f\ x \leq f\ y$

$(a::('a::\{\text{ord}, \text{zero}\})\ \text{matrix}) \leq b$

shows

$\text{apply-matrix}\ f\ a \leq \text{apply-matrix}\ f\ b$

by (*simp!* *add*: *le-matrix-def*)

lemma *le-combine-matrix*:

assumes

$f\ 0\ 0 = 0$

$! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow f\ a\ c \leq f\ b\ d$

$A \leq B$

$C \leq D$

shows

$\text{combine-matrix}\ f\ A\ C \leq \text{combine-matrix}\ f\ B\ D$

by (*simp!* *add*: *le-matrix-def*)

lemma *le-left-combine-matrix*:

assumes

$f\ 0\ 0 = 0$

$! a\ b\ c. a \leq b \longrightarrow f\ c\ a \leq f\ c\ b$

$A \leq B$

shows

$\text{combine-matrix}\ f\ C\ A \leq \text{combine-matrix}\ f\ C\ B$

by (*simp!* *add*: *le-matrix-def*)

lemma *le-right-combine-matrix*:

assumes

$f\ 0\ 0 = 0$

$! a\ b\ c. a \leq b \longrightarrow f\ a\ c \leq f\ b\ c$

$A \leq B$

shows

$\text{combine-matrix}\ f\ A\ C \leq \text{combine-matrix}\ f\ B\ C$

by (*simp!* *add*: *le-matrix-def*)

lemma *le-transpose-matrix*: $(A \leq B) = (\text{transpose-matrix}\ A \leq \text{transpose-matrix}\ B)$

by (*simp* *add*: *le-matrix-def*, *auto*)

lemma *le-foldseq*:

assumes

$! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow f\ a\ c \leq f\ b\ d$

$! i. i \leq n \longrightarrow s\ i \leq t\ i$

shows

$\text{foldseq}\ f\ s\ n \leq \text{foldseq}\ f\ t\ n$

proof –

have ! $s\ t.$ (! $i. i \leq n \longrightarrow s\ i \leq t\ i$) \longrightarrow $\text{foldseq}\ f\ s\ n \leq \text{foldseq}\ f\ t\ n$ **by**
(induct-tac n, simp-all!)
then show $\text{foldseq}\ f\ s\ n \leq \text{foldseq}\ f\ t\ n$ **by** *(simp!)*
qed

lemma *le-left-mult:*

assumes
! $a\ b\ c\ d.$ $a \leq b \ \& \ c \leq d \longrightarrow \text{fadd}\ a\ c \leq \text{fadd}\ b\ d$
! $c\ a\ b.$ $0 \leq c \ \& \ a \leq b \longrightarrow \text{fmul}\ c\ a \leq \text{fmul}\ c\ b$
! $a.$ $\text{fmul}\ 0\ a = 0$
! $a.$ $\text{fmul}\ a\ 0 = 0$
 $\text{fadd}\ 0\ 0 = 0$
 $0 \leq C$
 $A \leq B$
shows
 $\text{mult-matrix}\ \text{fmul}\ \text{fadd}\ C\ A \leq \text{mult-matrix}\ \text{fmul}\ \text{fadd}\ C\ B$
apply *(simp! add: le-matrix-def Rep-mult-matrix)*
apply *(auto)*
apply *(simplesubst foldseq-zerotail[of - - - max (ncols C) (max (nrows A) (nrows B))], simp-all add: nrows ncols max1 max2)+*
apply *(rule le-foldseq)*
by *(auto)*

lemma *le-right-mult:*

assumes
! $a\ b\ c\ d.$ $a \leq b \ \& \ c \leq d \longrightarrow \text{fadd}\ a\ c \leq \text{fadd}\ b\ d$
! $c\ a\ b.$ $0 \leq c \ \& \ a \leq b \longrightarrow \text{fmul}\ a\ c \leq \text{fmul}\ b\ c$
! $a.$ $\text{fmul}\ 0\ a = 0$
! $a.$ $\text{fmul}\ a\ 0 = 0$
 $\text{fadd}\ 0\ 0 = 0$
 $0 \leq C$
 $A \leq B$
shows
 $\text{mult-matrix}\ \text{fmul}\ \text{fadd}\ A\ C \leq \text{mult-matrix}\ \text{fmul}\ \text{fadd}\ B\ C$
apply *(simp! add: le-matrix-def Rep-mult-matrix)*
apply *(auto)*
apply *(simplesubst foldseq-zerotail[of - - - max (nrows C) (max (ncols A) (ncols B))], simp-all add: nrows ncols max1 max2)+*
apply *(rule le-foldseq)*
by *(auto)*

lemma *spec2:* ! $j\ i.$ $P\ j\ i \Longrightarrow P\ j\ i$ **by** *blast*

lemma *neg-imp:* $(\neg Q \longrightarrow \neg P) \Longrightarrow P \longrightarrow Q$ **by** *blast*

lemma *singleton-matrix-le[simp]:* $(\text{singleton-matrix}\ j\ i\ a \leq \text{singleton-matrix}\ j\ i\ b) = (a \leq (b::\text{order}))$
by *(auto simp add: le-matrix-def)*

lemma *singleton-le-zero[simp]:* $(\text{singleton-matrix}\ j\ i\ x \leq 0) = (x \leq (0::'a::\{\text{order},\text{zero}\}))$

```

apply (auto)
apply (simp add: le-matrix-def)
apply (drule-tac j=j and i=i in spec2)
apply (simp)
apply (simp add: le-matrix-def)
done

lemma singleton-ge-zero[simp]: ( $0 \leq \text{singleton-matrix } j \ i \ x = ((0::'a::\{\text{order},\text{zero}\}) \leq x)$ )
apply (auto)
apply (simp add: le-matrix-def)
apply (drule-tac j=j and i=i in spec2)
apply (simp)
apply (simp add: le-matrix-def)
done

lemma move-matrix-le-zero[simp]:  $0 \leq j \implies 0 \leq i \implies (\text{move-matrix } A \ j \ i \leq 0) = (A \leq (0::('a::\{\text{order},\text{zero}\}) \text{ matrix}))$ 
apply (auto simp add: le-matrix-def neg-def)
apply (drule-tac j=ja+(nat j) and i=ia+(nat i) in spec2)
apply (auto)
done

lemma move-matrix-zero-le[simp]:  $0 \leq j \implies 0 \leq i \implies (0 \leq \text{move-matrix } A \ j \ i) = ((0::('a::\{\text{order},\text{zero}\}) \text{ matrix}) \leq A)$ 
apply (auto simp add: le-matrix-def neg-def)
apply (drule-tac j=ja+(nat j) and i=ia+(nat i) in spec2)
apply (auto)
done

lemma move-matrix-le-move-matrix-iff[simp]:  $0 \leq j \implies 0 \leq i \implies (\text{move-matrix } A \ j \ i \leq \text{move-matrix } B \ j \ i) = (A \leq (B::('a::\{\text{order},\text{zero}\}) \text{ matrix}))$ 
apply (auto simp add: le-matrix-def neg-def)
apply (drule-tac j=ja+(nat j) and i=ia+(nat i) in spec2)
apply (auto)
done

end

theory Matrix=MatrixGeneral:

instance matrix :: (minus) minus
by intro-classes

instance matrix :: (plus) plus
by (intro-classes)

```

```

instance matrix :: ({plus,times}) times
by (intro-classes)

defs (overloaded)
  plus-matrix-def:  $A + B == combine-matrix (op +) A B$ 
  diff-matrix-def:  $A - B == combine-matrix (op -) A B$ 
  minus-matrix-def:  $- A == apply-matrix uminus A$ 
  times-matrix-def:  $A * B == mult-matrix (op *) (op +) A B$ 

lemma is-meet-combine-matrix-meet: is-meet (combine-matrix meet)
  by (simp-all add: is-meet-def le-matrix-def meet-left-le meet-right-le meet-imp-le)

lemma is-join-combine-matrix-join: is-join (combine-matrix join)
  by (simp-all add: is-join-def le-matrix-def join-left-le join-right-le join-imp-le)

instance matrix :: (lordered-ab-group) lordered-ab-group-meet
proof
  fix A B C :: ('a::lordered-ab-group) matrix
  show  $A + B + C = A + (B + C)$ 
    apply (simp add: plus-matrix-def)
    apply (rule combine-matrix-assoc[simplified associative-def, THEN spec, THEN spec, THEN spec])
    apply (simp-all add: add-assoc)
    done
  show  $A + B = B + A$ 
    apply (simp add: plus-matrix-def)
    apply (rule combine-matrix-commute[simplified commutative-def, THEN spec, THEN spec])
    apply (simp-all add: add-commute)
    done
  show  $0 + A = A$ 
    apply (simp add: plus-matrix-def)
    apply (rule combine-matrix-zero-l-neutral[simplified zero-l-neutral-def, THEN spec])
    apply (simp)
    done
  show  $- A + A = 0$ 
    by (simp add: plus-matrix-def minus-matrix-def Rep-matrix-inject[symmetric] ext)
  show  $A - B = A + - B$ 
    by (simp add: plus-matrix-def diff-matrix-def minus-matrix-def Rep-matrix-inject[symmetric] ext)
  show  $\exists m::'a matrix \Rightarrow 'a matrix \Rightarrow 'a matrix. is-meet m$ 
    by (auto intro: is-meet-combine-matrix-meet)
  assume A <= B
  then show C + A <= C + B
    apply (simp add: plus-matrix-def)
    apply (rule le-left-combine-matrix)
    apply (simp-all)

```

```

    done
  qed

  defs (overloaded)
    abs-matrix-def: abs (A::('a::lordered-ab-group) matrix) == join A (- A)

  instance matrix :: (lordered-ring) lordered-ring
  proof
    fix A B C :: ('a :: lordered-ring) matrix
    show A * B * C = A * (B * C)
      apply (simp add: times-matrix-def)
      apply (rule mult-matrix-assoc)
      apply (simp-all add: associative-def ring-eq-simps)
    done
    show (A + B) * C = A * C + B * C
      apply (simp add: times-matrix-def plus-matrix-def)
      apply (rule l-distributive-matrix[simplified l-distributive-def, THEN spec, THEN
spec, THEN spec])
      apply (simp-all add: associative-def commutative-def ring-eq-simps)
    done
    show A * (B + C) = A * B + A * C
      apply (simp add: times-matrix-def plus-matrix-def)
      apply (rule r-distributive-matrix[simplified r-distributive-def, THEN spec, THEN
spec, THEN spec])
      apply (simp-all add: associative-def commutative-def ring-eq-simps)
    done
    show abs A = join A (-A)
      by (simp add: abs-matrix-def)
    assume a: A ≤ B
    assume b: 0 ≤ C
    from a b show C * A ≤ C * B
      apply (simp add: times-matrix-def)
      apply (rule le-left-mult)
      apply (simp-all add: add-mono mult-left-mono)
    done
    from a b show A * C ≤ B * C
      apply (simp add: times-matrix-def)
      apply (rule le-right-mult)
      apply (simp-all add: add-mono mult-right-mono)
    done
  qed

  lemma Rep-matrix-add[simp]: Rep-matrix ((a::('a::lordered-ab-group)matrix)+b)
  j i = (Rep-matrix a j i) + (Rep-matrix b j i)
  by (simp add: plus-matrix-def)

  lemma Rep-matrix-mult: Rep-matrix ((a::('a::lordered-ring) matrix) * b) j i =
  foldseq (op +) (% k. (Rep-matrix a j k) * (Rep-matrix b k i)) (max (ncols a)
(nrows b))

```

```

apply (simp add: times-matrix-def)
apply (simp add: Rep-mult-matrix)
done

```

```

lemma apply-matrix-add: ! x y. f (x+y) = (f x) + (f y)  $\implies$  f 0 = (0::'a)  $\implies$ 
  apply-matrix f ((a::('a::lordered-ab-group) matrix) + b) = (apply-matrix f a) +
  (apply-matrix f b)
apply (subst Rep-matrix-inject[symmetric])
apply (rule ext)+
apply (simp)
done

```

```

lemma singleton-matrix-add: singleton-matrix j i ((a:::lordered-ab-group)+b) =
  (singleton-matrix j i a) + (singleton-matrix j i b)
apply (subst Rep-matrix-inject[symmetric])
apply (rule ext)+
apply (simp)
done

```

```

lemma nrows-mult: nrows ((A::('a::lordered-ring) matrix) * B) <= nrows A
by (simp add: times-matrix-def mult-nrows)

```

```

lemma ncols-mult: ncols ((A::('a::lordered-ring) matrix) * B) <= ncols B
by (simp add: times-matrix-def mult-ncols)

```

constdefs

```

  one-matrix :: nat  $\implies$  ('a::{zero,one}) matrix
  one-matrix n == Abs-matrix (% j i. if j = i & j < n then 1 else 0)

```

```

lemma Rep-one-matrix[simp]: Rep-matrix (one-matrix n) j i = (if (j = i & j <
  n) then 1 else 0)
apply (simp add: one-matrix-def)
apply (simpsubst RepAbs-matrix)
apply (rule exI[of - n], simp add: split-if)+
by (simp add: split-if)

```

```

lemma nrows-one-matrix[simp]: nrows ((one-matrix n) :: ('a::aclass-0-neq-1)matrix)
  = n (is ?r = -)

```

```

proof -
  have ?r <= n by (simp add: nrows-le)
  moreover have n <= ?r by (simp add: le-nrows, arith)
  ultimately show ?r = n by simp

```

qed

```

lemma ncols-one-matrix[simp]: ncols ((one-matrix n) :: ('a::aclass-0-neq-1)matrix)
  = n (is ?r = -)

```

```

proof -
  have ?r <= n by (simp add: ncols-le)

```

moreover have $n \leq r$ **by** (*simp add: le-ncols, arith*)
ultimately show $r = n$ **by** *simp*
qed

lemma *one-matrix-mult-right*[*simp*]: $\text{ncols } A \leq n \implies (A::('a::\{\text{lordered-ring, ring-1}\}) \text{matrix}) * (\text{one-matrix } n) = A$
apply (*subst Rep-matrix-inject*[*THEN sym*])
apply (*rule ext*)
apply (*simp add: times-matrix-def Rep-mult-matrix*)
apply (*rule-tac j1=xa in ssubst*[*OF foldseq-almostzero*])
apply (*simp-all*)
by (*simp add: max-def ncols*)

lemma *one-matrix-mult-left*[*simp*]: $\text{nrows } A \leq n \implies (\text{one-matrix } n) * A = (A::('a::\{\text{lordered-ring, ring-1}\}) \text{matrix})$
apply (*subst Rep-matrix-inject*[*THEN sym*])
apply (*rule ext*)
apply (*simp add: times-matrix-def Rep-mult-matrix*)
apply (*rule-tac j1=x in ssubst*[*OF foldseq-almostzero*])
apply (*simp-all*)
by (*simp add: max-def nrows*)

lemma *transpose-matrix-mult*: $\text{transpose-matrix } ((A::('a::\{\text{lordered-ring, comm-ring}\}) \text{matrix}) * B) = (\text{transpose-matrix } B) * (\text{transpose-matrix } A)$
apply (*simp add: times-matrix-def*)
apply (*subst transpose-mult-matrix*)
apply (*simp-all add: mult-commute*)
done

lemma *transpose-matrix-add*: $\text{transpose-matrix } ((A::('a::\text{lordered-ab-group}) \text{matrix}) + B) = \text{transpose-matrix } A + \text{transpose-matrix } B$
by (*simp add: plus-matrix-def transpose-combine-matrix*)

lemma *transpose-matrix-diff*: $\text{transpose-matrix } ((A::('a::\text{lordered-ab-group}) \text{matrix}) - B) = \text{transpose-matrix } A - \text{transpose-matrix } B$
by (*simp add: diff-matrix-def transpose-combine-matrix*)

lemma *transpose-matrix-minus*: $\text{transpose-matrix } (- (A::('a::\text{lordered-ring}) \text{matrix})) = - \text{transpose-matrix } (A::('a::\text{lordered-ring}) \text{matrix})$
by (*simp add: minus-matrix-def transpose-apply-matrix*)

constdefs

right-inverse-matrix :: $('a::\{\text{lordered-ring, ring-1}\}) \text{matrix} \Rightarrow 'a \text{matrix} \Rightarrow \text{bool}$
right-inverse-matrix $A X == (A * X = \text{one-matrix } (\text{max } (\text{nrows } A) (\text{ncols } X))) \wedge \text{nrows } X \leq \text{ncols } A$
left-inverse-matrix :: $('a::\{\text{lordered-ring, ring-1}\}) \text{matrix} \Rightarrow 'a \text{matrix} \Rightarrow \text{bool}$
left-inverse-matrix $A X == (X * A = \text{one-matrix } (\text{max } (\text{nrows } X) (\text{ncols } A))) \wedge \text{ncols } X \leq \text{nrows } A$
inverse-matrix :: $('a::\{\text{lordered-ring, ring-1}\}) \text{matrix} \Rightarrow 'a \text{matrix} \Rightarrow \text{bool}$

inverse-matrix A X == (right-inverse-matrix A X) ∧ (left-inverse-matrix A X)

lemma *right-inverse-matrix-dim: right-inverse-matrix A X ==> nrows A = ncols X*
apply (*insert ncols-mult[of A X], insert nrows-mult[of A X]*)
by (*simp add: right-inverse-matrix-def*)

lemma *left-inverse-matrix-dim: left-inverse-matrix A Y ==> ncols A = nrows Y*
apply (*insert ncols-mult[of Y A], insert nrows-mult[of Y A]*)
by (*simp add: left-inverse-matrix-def*)

lemma *left-right-inverse-matrix-unique:*
assumes *left-inverse-matrix A Y right-inverse-matrix A X*
shows *X = Y*
proof –
have *Y = Y * one-matrix (nrows A)*
apply (*subst one-matrix-mult-right*)
apply (*insert prems*)
by (*simp-all add: left-inverse-matrix-def*)
also have *... = Y * (A * X)*
apply (*insert prems*)
apply (*frule right-inverse-matrix-dim*)
by (*simp add: right-inverse-matrix-def*)
also have *... = (Y * A) * X* **by** (*simp add: mult-assoc*)
also have *... = X*
apply (*insert prems*)
apply (*frule left-inverse-matrix-dim*)
apply (*simp-all add: left-inverse-matrix-def right-inverse-matrix-def one-matrix-mult-left*)
done
ultimately show *X = Y* **by** (*simp*)
qed

lemma *inverse-matrix-inject: [inverse-matrix A X; inverse-matrix A Y] ==> X = Y*
by (*auto simp add: inverse-matrix-def left-right-inverse-matrix-unique*)

lemma *one-matrix-inverse: inverse-matrix (one-matrix n) (one-matrix n)*
by (*simp add: inverse-matrix-def left-inverse-matrix-def right-inverse-matrix-def*)

lemma *zero-imp-mult-zero: (a::'a::ring) = 0 | b = 0 ==> a * b = 0*
by *auto*

lemma *Rep-matrix-zero-imp-mult-zero:*
! *j i k. (Rep-matrix A j k = 0) | (Rep-matrix B k i) = 0 ==> A * B = (0::('a::lordered-ring) matrix)*
apply (*subst Rep-matrix-inject[symmetric]*)
apply (*rule ext*)
apply (*auto simp add: Rep-matrix-mult foldseq-zero zero-imp-mult-zero*)
done

```

lemma add-nrows: nrows (A::('a::comm-monoid-add) matrix) <= u  $\implies$  nrows B
<= u  $\implies$  nrows (A + B) <= u
apply (simp add: plus-matrix-def)
apply (rule combine-nrows)
apply (simp-all)
done

```

```

lemma move-matrix-row-mult: move-matrix ((A::('a::lordered-ring) matrix) * B)
j 0 = (move-matrix A j 0) * B
apply (subst Rep-matrix-inject[symmetric])
apply (rule ext)+
apply (auto simp add: Rep-matrix-mult foldseq-zero)
apply (rule-tac foldseq-zerotail[symmetric])
apply (auto simp add: nrows zero-imp-mult-zero max2)
apply (rule order-trans)
apply (rule ncols-move-matrix-le)
apply (simp add: max1)
done

```

```

lemma move-matrix-col-mult: move-matrix ((A::('a::lordered-ring) matrix) * B)
0 i = A * (move-matrix B 0 i)
apply (subst Rep-matrix-inject[symmetric])
apply (rule ext)+
apply (auto simp add: Rep-matrix-mult foldseq-zero)
apply (rule-tac foldseq-zerotail[symmetric])
apply (auto simp add: ncols zero-imp-mult-zero max1)
apply (rule order-trans)
apply (rule nrows-move-matrix-le)
apply (simp add: max2)
done

```

```

lemma move-matrix-add: ((move-matrix (A + B) j i)::('a::lordered-ab-group)
matrix) = (move-matrix A j i) + (move-matrix B j i)
apply (subst Rep-matrix-inject[symmetric])
apply (rule ext)+
apply (simp)
done

```

```

lemma move-matrix-mult: move-matrix ((A::('a::lordered-ring) matrix)*B) j i =
(move-matrix A j 0) * (move-matrix B 0 i)
by (simp add: move-matrix-ortho[of A*B] move-matrix-col-mult move-matrix-row-mult)

```

constdefs

```

scalar-mult :: ('a::lordered-ring)  $\Rightarrow$  'a matrix  $\Rightarrow$  'a matrix
scalar-mult a m == apply-matrix (op * a) m

```

```

lemma scalar-mult-zero[simp]: scalar-mult y 0 = 0
by (simp add: scalar-mult-def)

```

```

lemma scalar-mult-add: scalar-mult  $y$  ( $a+b$ ) = (scalar-mult  $y$   $a$ ) + (scalar-mult  $y$ 
b)
  by (simp add: scalar-mult-def apply-matrix-add ring-eq-simps)

lemma Rep-scalar-mult[simp]: Rep-matrix (scalar-mult  $y$   $a$ )  $j$   $i$  =  $y$  * (Rep-matrix
 $a$   $j$   $i$ )
  by (simp add: scalar-mult-def)

lemma scalar-mult-singleton[simp]: scalar-mult  $y$  (singleton-matrix  $j$   $i$   $x$ ) = singleton-matrix
 $j$   $i$  ( $y$  *  $x$ )
  apply (subst Rep-matrix-inject[symmetric])
  apply (rule ext)+
  apply (auto)
  done

lemma Rep-minus[simp]: Rep-matrix ( $-(A:::\text{ordered-ab-group})$ )  $x$   $y$  =  $-$  (Rep-matrix
 $A$   $x$   $y$ )
  by (simp add: minus-matrix-def)

lemma join-matrix: join ( $A::('a::\text{ordered-ring})$  matrix)  $B$  = combine-matrix join
 $A$   $B$ 
  apply (insert join-unique[of (combine-matrix join)::('a matrix  $\Rightarrow$  'a matrix  $\Rightarrow$ 
' a matrix), simplified is-join-combine-matrix-join])
  apply (simp)
  done

lemma meet-matrix: meet ( $A::('a::\text{ordered-ring})$  matrix)  $B$  = combine-matrix
meet  $A$   $B$ 
  apply (insert meet-unique[of (combine-matrix meet)::('a matrix  $\Rightarrow$  'a matrix  $\Rightarrow$ 
' a matrix), simplified is-meet-combine-matrix-meet])
  apply (simp)
  done

lemma Rep-abs[simp]: Rep-matrix (abs ( $A:::\text{ordered-ring}$ ))  $x$   $y$  = abs (Rep-matrix
 $A$   $x$   $y$ )
  by (simp add: abs-lattice join-matrix)

end

```

```

theory SparseMatrix imports Matrix begin

```

```

types

```

```

  'a svec = ( $\text{nat}$  * 'a) list
  'a spmat = ('a svec) svec

```

```

consts

```

sparse-row-vector :: ('a::lordered-ring) spvec \Rightarrow 'a matrix
sparse-row-matrix :: ('a::lordered-ring) spmat \Rightarrow 'a matrix

defs

sparse-row-vector-def : *sparse-row-vector* arr == foldl (% m x. m + (singleton-matrix 0 (fst x) (snd x))) 0 arr

sparse-row-matrix-def : *sparse-row-matrix* arr == foldl (% m r. m + (move-matrix (sparse-row-vector (snd r)) (int (fst r)) 0)) 0 arr

lemma *sparse-row-vector-empty[simp]*: *sparse-row-vector* [] = 0
by (*simp add: sparse-row-vector-def*)

lemma *sparse-row-matrix-empty[simp]*: *sparse-row-matrix* [] = 0
by (*simp add: sparse-row-matrix-def*)

lemma *foldl-distrstart[rule-format]*: ! a x y. (f (g x y) a = g x (f y a)) \implies ! x y.
(foldl f (g x y) l = g x (foldl f y l))
by (*induct l, auto*)

lemma *sparse-row-vector-cons[simp]*: *sparse-row-vector* (a#arr) = (singleton-matrix 0 (fst a) (snd a)) + (sparse-row-vector arr)
apply (*induct arr*)
apply (*auto simp add: sparse-row-vector-def*)
apply (*simp add: foldl-distrstart[of $\lambda m x. m +$ singleton-matrix 0 (fst x) (snd x) $\lambda x m. singleton-matrix 0 (fst x) (snd x) + m]$)*)
done

lemma *sparse-row-vector-append[simp]*: *sparse-row-vector* (a @ b) = (sparse-row-vector a) + (sparse-row-vector b)
by (*induct a, auto*)

lemma *nrows-spvec[simp]*: *nrows* (sparse-row-vector x) <= (Suc 0)
apply (*induct x*)
apply (*simp-all add: add-nrows*)
done

lemma *sparse-row-matrix-cons*: *sparse-row-matrix* (a#arr) = ((move-matrix (sparse-row-vector (snd a)) (int (fst a)) 0)) + *sparse-row-matrix* arr
apply (*induct arr*)
apply (*auto simp add: sparse-row-matrix-def*)
apply (*simp add: foldl-distrstart[of $\lambda m x. m +$ (move-matrix (sparse-row-vector (snd x)) (int (fst x)) 0) $\% a m. (move-matrix (sparse-row-vector (snd a)) (int (fst a)) 0) + m]$)*)
done

lemma *sparse-row-matrix-append*: *sparse-row-matrix* (arr@brr) = (sparse-row-matrix arr) + (sparse-row-matrix brr)
apply (*induct arr*)
apply (*auto simp add: sparse-row-matrix-cons*)

```

done

consts
  sorted-spvec :: 'a spvec ⇒ bool
  sorted-spmat :: 'a spmat ⇒ bool

primrec
  sorted-spmat [] = True
  sorted-spmat (a#as) = ((sorted-spvec (snd a)) & (sorted-spmat as))

primrec
  sorted-spvec [] = True
  sorted-spvec-step: sorted-spvec (a#as) = (case as of [] ⇒ True | b#bs ⇒ ((fst a
  < fst b) & (sorted-spvec as)))

declare sorted-spvec.simps [simp del]

lemma sorted-spvec-empty[simp]: sorted-spvec [] = True
by (simp add: sorted-spvec.simps)

lemma sorted-spvec-cons1: sorted-spvec (a#as) ⇒ sorted-spvec as
apply (induct as)
apply (auto simp add: sorted-spvec.simps)
done

lemma sorted-spvec-cons2: sorted-spvec (a#b#t) ⇒ sorted-spvec (a#t)
apply (induct t)
apply (auto simp add: sorted-spvec.simps)
done

lemma sorted-spvec-cons3: sorted-spvec(a#b#t) ⇒ fst a < fst b
apply (auto simp add: sorted-spvec.simps)
done

lemma sorted-sparse-row-vector-zero[rule-format]: m ≤ n → sorted-spvec ((n,a)#arr)
→ Rep-matrix (sparse-row-vector arr) j m = 0
apply (induct arr)
apply (auto)
apply (frule sorted-spvec-cons2, simp)+
apply (frule sorted-spvec-cons3, simp)
done

lemma sorted-sparse-row-matrix-zero[rule-format]: m ≤ n → sorted-spvec ((n,a)#arr)
→ Rep-matrix (sparse-row-matrix arr) m j = 0
apply (induct arr)
apply (auto)
apply (frule sorted-spvec-cons2, simp)
apply (frule sorted-spvec-cons3, simp)
apply (simp add: sparse-row-matrix-cons neg-def)

```

done

consts

$abs\text{-}spvec :: ('a::lordered\text{-}ring) \text{ } spvec \Rightarrow 'a \text{ } spvec$
 $minus\text{-}spvec :: ('a::lordered\text{-}ring) \text{ } spvec \Rightarrow 'a \text{ } spvec$
 $smult\text{-}spvec :: ('a::lordered\text{-}ring) \Rightarrow 'a \text{ } spvec \Rightarrow 'a \text{ } spvec$
 $addmult\text{-}spvec :: ('a::lordered\text{-}ring) * 'a \text{ } spvec * 'a \text{ } spvec \Rightarrow 'a \text{ } spvec$

primrec

$minus\text{-}spvec \ [] = []$
 $minus\text{-}spvec (a\#as) = (fst\ a, -(snd\ a))\#(minus\text{-}spvec\ as)$

primrec

$abs\text{-}spvec \ [] = []$
 $abs\text{-}spvec (a\#as) = (fst\ a, abs\ (snd\ a))\#(abs\text{-}spvec\ as)$

lemma *sparse-row-vector-minus:*

$sparse\text{-}row\text{-}vector\ (minus\text{-}spvec\ v) = -\ (sparse\text{-}row\text{-}vector\ v)$
apply (*induct v*)
apply (*simp-all add: sparse-row-vector-cons*)
apply (*simp add: Rep-matrix-inject[symmetric]*)
apply (*rule ext*)
apply *simp*
done

lemma *sparse-row-vector-abs:*

$sorted\text{-}spvec\ v \Longrightarrow sparse\text{-}row\text{-}vector\ (abs\text{-}spvec\ v) = abs\ (sparse\text{-}row\text{-}vector\ v)$
apply (*induct v*)
apply (*simp-all add: sparse-row-vector-cons*)
apply (*frule-tac sorted-spvec-cons1, simp*)
apply (*simp only: Rep-matrix-inject[symmetric]*)
apply (*rule ext*)
apply *auto*
apply (*subgoal-tac Rep-matrix (sparse-row-vector v) 0 a = 0*)
apply (*simp*)
apply (*rule sorted-sparse-row-vector-zero*)
apply *auto*
done

lemma *sorted-spvec-minus-spvec:*

$sorted\text{-}spvec\ v \Longrightarrow sorted\text{-}spvec\ (minus\text{-}spvec\ v)$
apply (*induct v*)
apply (*simp*)
apply (*frule sorted-spvec-cons1, simp*)
apply (*simp add: sorted-spvec.simps split:list.split-asm*)
done

lemma *sorted-spvec-minus-spvec:*

$sorted\text{-}spvec\ v \Longrightarrow sorted\text{-}spvec\ (minus\text{-}spvec\ v)$

```

apply (induct v)
apply (simp)
apply (frule sorted-spvec-cons1, simp)
apply (simp add: sorted-spvec.simps split:list.split-asm)
done

lemma sorted-spvec-abs-spvec:
  sorted-spvec v  $\implies$  sorted-spvec (abs-spvec v)
apply (induct v)
apply (simp)
apply (frule sorted-spvec-cons1, simp)
apply (simp add: sorted-spvec.simps split:list.split-asm)
done

defs
  smult-spvec-def: smult-spvec y arr == map (% a. (fst a, y * snd a)) arr

lemma smult-spvec-empty[simp]: smult-spvec y [] = []
by (simp add: smult-spvec-def)

lemma smult-spvec-cons: smult-spvec y (a#arr) = (fst a, y * (snd a)) # (smult-spvec y arr)
by (simp add: smult-spvec-def)

recdef addmult-spvec measure (% (y, a, b). length a + (length b))
  addmult-spvec (y, arr, []) = arr
  addmult-spvec (y, [], brr) = smult-spvec y brr
  addmult-spvec (y, a#arr, b#brr) = (
    if (fst a) < (fst b) then (a#(addmult-spvec (y, arr, b#brr)))
    else (if (fst b < fst a) then ((fst b, y * (snd b))#(addmult-spvec (y, a#arr, brr)))
    else ((fst a, (snd a)+ y*(snd b))#(addmult-spvec (y, arr,brr)))))

lemma addmult-spvec-empty1[simp]: addmult-spvec (y, [], a) = smult-spvec y a
by (induct a, auto)

lemma addmult-spvec-empty2[simp]: addmult-spvec (y, a, []) = a
by (induct a, auto)

lemma sparse-row-vector-map: (! x y. f (x+y) = (f x) + (f y))  $\implies$  (f::'a $\Rightarrow$ ('a::lordered-ring))
0 = 0  $\implies$ 
  sparse-row-vector (map (% x. (fst x, f (snd x))) a) = apply-matrix f (sparse-row-vector a)
apply (induct a)
apply (simp-all add: apply-matrix-add)
done

lemma sparse-row-vector-smult: sparse-row-vector (smult-spvec y a) = scalar-mult y (sparse-row-vector a)

```

```

apply (induct a)
apply (simp-all add: smult-spvec-cons scalar-mult-add)
done

```

```

lemma sparse-row-vector-addmult-spvec: sparse-row-vector (addmult-spvec (y::'a::lordered-ring,
a, b)) =
  (sparse-row-vector a) + (scalar-mult y (sparse-row-vector b))
apply (rule addmult-spvec.induct[of - y])
apply (simp add: scalar-mult-add smult-spvec-cons sparse-row-vector-smult singleton-matrix-add)+
done

```

```

lemma sorted-smult-spvec[rule-format]: sorted-spvec a  $\implies$  sorted-spvec (smult-spvec
y a)
apply (auto simp add: smult-spvec-def)
apply (induct a)
apply (auto simp add: sorted-spvec.simps split:list.split-asm)
done

```

```

lemma sorted-spvec-addmult-spvec-helper:  $\llbracket$ sorted-spvec (addmult-spvec (y, (a, b)
# arr, brr)); aa < a; sorted-spvec ((a, b) # arr);
  sorted-spvec ((aa, ba) # brr) $\rrbracket \implies$  sorted-spvec ((aa, y * ba) # addmult-spvec
(y, (a, b) # arr, brr))
apply (induct brr)
apply (auto simp add: sorted-spvec.simps)
apply (simp split: list.split)
apply (auto)
apply (simp split: list.split)
apply (auto)
done

```

```

lemma sorted-spvec-addmult-spvec-helper2:
 $\llbracket$ sorted-spvec (addmult-spvec (y, arr, (aa, ba) # brr)); a < aa; sorted-spvec ((a,
b) # arr); sorted-spvec ((aa, ba) # brr) $\rrbracket$ 
 $\implies$  sorted-spvec ((a, b) # addmult-spvec (y, arr, (aa, ba) # brr))
apply (induct arr)
apply (auto simp add: smult-spvec-def sorted-spvec.simps)
apply (simp split: list.split)
apply (auto)
done

```

```

lemma sorted-spvec-addmult-spvec-helper3[rule-format]:
  sorted-spvec (addmult-spvec (y, arr, brr))  $\longrightarrow$  sorted-spvec ((aa, b) # arr)  $\longrightarrow$ 
  sorted-spvec ((aa, ba) # brr)
 $\longrightarrow$  sorted-spvec ((aa, b + y * ba) # (addmult-spvec (y, arr, brr)))
apply (rule addmult-spvec.induct[of - y arr brr])
apply (simp-all add: sorted-spvec.simps smult-spvec-def)
done

```

```

lemma sorted-addmult-spvec[rule-format]: sorted-spvec a  $\longrightarrow$  sorted-spvec b  $\longrightarrow$ 

```

```

sorted-spvec (addmult-spvec (y, a, b))
  apply (rule addmult-spvec.induct[of - y a b])
  apply (simp-all add: sorted-smult-spvec)
  apply (rule conjI, intro strip)
  apply (case-tac ~ (a < aa))
  apply (simp-all)
  apply (frule-tac as=brr in sorted-spvec-cons1)
  apply (simp add: sorted-spvec-addmult-spvec-helper)
  apply (intro strip | rule conjI)+
  apply (frule-tac as=arr in sorted-spvec-cons1)
  apply (simp add: sorted-spvec-addmult-spvec-helper2)
  apply (intro strip)
  apply (frule-tac as=arr in sorted-spvec-cons1)
  apply (frule-tac as=brr in sorted-spvec-cons1)
  apply (simp)
  apply (simp-all add: sorted-spvec-addmult-spvec-helper3)
done

```

consts

```

mult-spvec-spmat :: ('a::lordered-ring) spvec * 'a spvec * 'a spmat  $\Rightarrow$  'a spvec

```

recdef *mult-spvec-spmat measure* (% (c, arr, brr). (length arr) + (length brr))

```

mult-spvec-spmat (c, [], brr) = c
mult-spvec-spmat (c, arr, []) = c
mult-spvec-spmat (c, a#arr, b#brr) = (
  if ((fst a) < (fst b)) then (mult-spvec-spmat (c, arr, b#brr))
  else (if ((fst b) < (fst a)) then (mult-spvec-spmat (c, a#arr, brr))
  else (mult-spvec-spmat (addmult-spvec (snd a, c, snd b), arr, brr)))

```

lemma *sparse-row-mult-spvec-spmat*[rule-format]: sorted-spvec (a::('a::lordered-ring) spvec) \longrightarrow sorted-spvec B \longrightarrow

```

sparse-row-vector (mult-spvec-spmat (c, a, B)) = (sparse-row-vector c) + (sparse-row-vector a) * (sparse-row-matrix B)

```

proof –

```

have comp-1: !! a b. a < b  $\implies$  Suc 0 <= nat ((int b)–(int a)) by arith

```

```

have not-iff: !! a b. a = b  $\implies$  (~ a) = (~ b) by simp

```

```

have max-helper: !! a b. ~ (a <= max (Suc a) b)  $\implies$  False

```

```

  by arith

```

```

{

```

```
  fix a

```

```
  fix v

```

```
  assume a:a < nrows(sparse-row-vector v)

```

```
  have b:nrows(sparse-row-vector v) <= 1 by simp

```

```
  note dummy = less-le-trans[of a nrows (sparse-row-vector v) 1, OF a b]

```

```
  then have a = 0 by simp

```

```

}

```

```

note nrows-helper = this

```

```

show ?thesis

```

```
  apply (rule mult-spvec-spmat.induct)

```

```

apply simp+
apply (rule conjI)
apply (intro strip)
apply (frule-tac as=brr in sorted-spvec-cons1)
apply (simp add: ring-eq-simps sparse-row-matrix-cons)
apply (simplesubst Rep-matrix-zero-imp-mult-zero)
apply (simp)
apply (intro strip)
apply (rule disjI2)
apply (intro strip)
apply (subst nrows)
apply (rule order-trans[of - 1])
apply (simp add: comp-1)+
apply (subst Rep-matrix-zero-imp-mult-zero)
apply (intro strip)
apply (case-tac k <= aa)
apply (rule-tac m1 = k and n1 = a and a1 = b in ssubst[OF sorted-sparse-row-vector-zero])
apply (simp-all)
apply (rule impI)
apply (rule disjI2)
apply (rule nrows)
apply (rule order-trans[of - 1])
apply (simp-all add: comp-1)

apply (intro strip | rule conjI)+
apply (frule-tac as=arr in sorted-spvec-cons1)
apply (simp add: ring-eq-simps)
apply (subst Rep-matrix-zero-imp-mult-zero)
apply (simp)
apply (rule disjI2)
apply (intro strip)
apply (simp add: sparse-row-matrix-cons neg-def)
apply (case-tac a <= aa)
apply (erule sorted-sparse-row-matrix-zero)
apply (simp-all)
apply (intro strip)
apply (case-tac a=aa)
apply (simp-all)
apply (frule-tac as=arr in sorted-spvec-cons1)
apply (frule-tac as=brr in sorted-spvec-cons1)
apply (simp add: sparse-row-matrix-cons ring-eq-simps sparse-row-vector-addmult-spvec)
apply (rule-tac B1 = sparse-row-matrix brr in ssubst[OF Rep-matrix-zero-imp-mult-zero])
apply (auto)
apply (rule sorted-sparse-row-matrix-zero)
apply (simp-all)
apply (rule-tac A1 = sparse-row-vector arr in ssubst[OF Rep-matrix-zero-imp-mult-zero])
apply (auto)
apply (rule-tac m=k and n = aa and a = b and arr=arr in sorted-sparse-row-vector-zero)
apply (simp-all)

```

```

apply (simp add: neg-def)
apply (drule nrows-notzero)
apply (drule nrows-helper)
apply (arith)

apply (subst Rep-matrix-inject[symmetric])
apply (rule ext)+
apply (simp)
apply (subst Rep-matrix-mult)
apply (rule-tac j1=aa in ssubst[OF foldseq-almostzero])
apply (simp-all)
apply (intro strip, rule conjI)
apply (intro strip)
apply (drule-tac max-helper)
apply (simp)
apply (auto)
apply (rule zero-imp-mult-zero)
apply (rule disjI2)
apply (rule nrows)
apply (rule order-trans[of - 1])
apply (simp)
apply (simp)
done

```

qed

```

lemma sorted-mult-spvec-spmat[rule-format]:
  sorted-spvec (c::('a::lordered-ring) spvec)  $\longrightarrow$  sorted-spmat B  $\longrightarrow$  sorted-spvec
  (mult-spvec-spmat (c, a, B))
apply (rule mult-spvec-spmat.induct[of - c a B])
apply (simp-all add: sorted-addmult-spvec)
done

```

consts

```

mult-spmat :: ('a::lordered-ring) spmat  $\Rightarrow$  'a spmat  $\Rightarrow$  'a spmat

```

primrec

```

mult-spmat [] A = []
mult-spmat (a#as) A = (fst a, mult-spvec-spmat ([], snd a, A))#(mult-spmat as A)

```

lemma *sparse-row-mult-spmat[rule-format]*:

```

sorted-spmat A  $\longrightarrow$  sorted-spvec B  $\longrightarrow$  sparse-row-matrix (mult-spmat A B) =
  (sparse-row-matrix A) * (sparse-row-matrix B)
apply (induct A)
apply (auto simp add: sparse-row-matrix-cons sparse-row-mult-spvec-spmat ring-eq-simps
  move-matrix-mult)
done

```

lemma *sorted-spvec-mult-spmat[rule-format]*:

```

sorted-spvec (A::('a::lordered-ring) spmat)  $\longrightarrow$  sorted-spvec (mult-spmat A B)
apply (induct A)
apply (auto)
apply (drule sorted-spvec-cons1, simp)
apply (case-tac A)
apply (auto simp add: sorted-spvec.simps)
done

lemma sorted-spmat-mult-spmat[rule-format]:
sorted-spmat (B::('a::lordered-ring) spmat)  $\longrightarrow$  sorted-spmat (mult-spmat A B)
apply (induct A)
apply (auto simp add: sorted-mult-spvec-spmat)
done

consts
add-spvec :: ('a::lordered-ab-group) spvec * 'a spvec  $\Rightarrow$  'a spvec
add-spmat :: ('a::lordered-ab-group) spmat * 'a spmat  $\Rightarrow$  'a spmat

recdef add-spvec measure (% (a, b). length a + (length b))
add-spvec (arr, []) = arr
add-spvec ([], brr) = brr
add-spvec (a#arr, b#brr) = (
  if (fst a) < (fst b) then (a#(add-spvec (arr, b#brr)))
  else (if (fst b < fst a) then (b#(add-spvec (a#arr, brr)))
  else ((fst a, (snd a)+(snd b))#(add-spvec (arr,brr))))

lemma add-spvec-empty1[simp]: add-spvec ([], a) = a
by (induct a, auto)

lemma add-spvec-empty2[simp]: add-spvec (a, []) = a
by (induct a, auto)

lemma sparse-row-vector-add: sparse-row-vector (add-spvec (a,b)) = (sparse-row-vector a) + (sparse-row-vector b)
apply (rule add-spvec.induct[of - a b])
apply (simp-all add: singleton-matrix-add)
done

recdef add-spmat measure (% (A,B). (length A)+(length B))
add-spmat ([], bs) = bs
add-spmat (as, []) = as
add-spmat (a#as, b#bs) = (
  if fst a < fst b then
    (a#(add-spmat (as, b#bs)))
  else (if fst b < fst a then
    (b#(add-spmat (a#as, bs)))
  else
    ((fst a, add-spvec (snd a, snd b))#(add-spmat (as, bs))))

```

lemma *sparse-row-add-spmat*: *sparse-row-matrix* (add-spmat (A, B)) = (*sparse-row-matrix* A) + (*sparse-row-matrix* B)
apply (rule add-spmat.induct)
apply (auto simp add: *sparse-row-matrix-cons sparse-row-vector-add move-matrix-add*)
done

lemma *sorted-add-spvec-helper1*[*rule-format*]: *add-spvec* ((a,b)#arr, brr) = (ab, bb) # list \longrightarrow (ab = a | (brr \neq [] & ab = fst (hd brr)))
proof -
have (! x ab a. x = (a,b)#arr \longrightarrow *add-spvec* (x, brr) = (ab, bb) # list \longrightarrow (ab = a | (ab = fst (hd brr))))
by (rule add-spvec.induct[of - - brr], auto)
then show ?thesis
by (case-tac brr, auto)
qed

lemma *sorted-add-spmat-helper1*[*rule-format*]: *add-spmat* ((a,b)#arr, brr) = (ab, bb) # list \longrightarrow (ab = a | (brr \neq [] & ab = fst (hd brr)))
proof -
have (! x ab a. x = (a,b)#arr \longrightarrow *add-spmat* (x, brr) = (ab, bb) # list \longrightarrow (ab = a | (ab = fst (hd brr))))
by (rule add-spmat.induct[of - - brr], auto)
then show ?thesis
by (case-tac brr, auto)
qed

lemma *sorted-add-spvec-helper*[*rule-format*]: *add-spvec* (arr, brr) = (ab, bb) # list \longrightarrow ((arr \neq [] & ab = fst (hd arr)) | (brr \neq [] & ab = fst (hd brr)))
apply (rule add-spvec.induct[of - arr brr])
apply (auto)
done

lemma *sorted-add-spmat-helper*[*rule-format*]: *add-spmat* (arr, brr) = (ab, bb) # list \longrightarrow ((arr \neq [] & ab = fst (hd arr)) | (brr \neq [] & ab = fst (hd brr)))
apply (rule add-spmat.induct[of - arr brr])
apply (auto)
done

lemma *add-spvec-commute*: *add-spvec* (a, b) = *add-spvec* (b, a)
by (rule add-spvec.induct[of - a b], auto)

lemma *add-spmat-commute*: *add-spmat* (a, b) = *add-spmat* (b, a)
apply (rule add-spmat.induct[of - a b])
apply (simp-all add: *add-spvec-commute*)
done

lemma *sorted-add-spvec-helper2*: *add-spvec* ((a,b)#arr, brr) = (ab, bb) # list \implies aa < a \implies *sorted-spvec* ((aa, ba) # brr) \implies aa < ab
apply (drule *sorted-add-spvec-helper1*)

```

apply (auto)
apply (case-tac brr)
apply (simp-all)
apply (drule-tac sorted-spvec-cons3)
apply (simp)
done

```

```

lemma sorted-add-spmat-helper2: add-spmat ((a,b)#arr, brr) = (ab, bb) # list
 $\implies aa < a \implies \text{sorted-spvec} ((aa, ba) \# brr) \implies aa < ab$ 
apply (drule sorted-add-spmat-helper1)
apply (auto)
apply (case-tac brr)
apply (simp-all)
apply (drule-tac sorted-spvec-cons3)
apply (simp)
done

```

```

lemma sorted-spvec-add-spvec[rule-format]: sorted-spvec a  $\longrightarrow$  sorted-spvec b  $\longrightarrow$ 
sorted-spvec (add-spvec (a, b))
apply (rule add-spvec.induct[of - a b])
apply (simp-all)
apply (rule conjI)
apply (intro strip)
apply (simp)
apply (frule-tac as=brr in sorted-spvec-cons1)
apply (simp)
apply (subst sorted-spvec-step)
apply (simp split: list.split)
apply (clarify, simp)
apply (simp add: sorted-add-spvec-helper2)
apply (clarify)
apply (rule conjI)
apply (case-tac a=aa)
apply (simp)
apply (clarify)
apply (frule-tac as=arr in sorted-spvec-cons1, simp)
apply (subst sorted-spvec-step)
apply (simp split: list.split)
apply (clarify, simp)
apply (simp add: sorted-add-spvec-helper2 add-spvec-commute)
apply (case-tac a=aa)
apply (simp-all)
apply (clarify)
apply (frule-tac as=arr in sorted-spvec-cons1)
apply (frule-tac as=brr in sorted-spvec-cons1)
apply (simp)
apply (subst sorted-spvec-step)
apply (simp split: list.split)
apply (clarify, simp)

```

```

apply (drule-tac sorted-add-spvec-helper)
apply (auto)
apply (case-tac arr)
apply (simp-all)
apply (drule sorted-spvec-cons3)
apply (simp)
apply (case-tac brr)
apply (simp-all)
apply (drule sorted-spvec-cons3)
apply (simp)
done

```

```

lemma sorted-spvec-add-spmat[rule-format]: sorted-spvec A  $\longrightarrow$  sorted-spvec B
 $\longrightarrow$  sorted-spvec (add-spmat (A, B))
apply (rule add-spmat.induct[of - A B])
apply (simp-all)
apply (rule conjI)
apply (intro strip)
apply (simp)
apply (frule-tac as=bs in sorted-spvec-cons1)
apply (simp)
apply (subst sorted-spvec-step)
apply (simp split: list.split)
apply (clarify, simp)
apply (simp add: sorted-add-spmat-helper2)
apply (clarify)
apply (rule conjI)
apply (case-tac a=aa)
apply (simp)
apply (clarify)
apply (frule-tac as=as in sorted-spvec-cons1, simp)
apply (subst sorted-spvec-step)
apply (simp split: list.split)
apply (clarify, simp)
apply (simp add: sorted-add-spmat-helper2 add-spmat-commute)
apply (case-tac a=aa)
apply (simp-all)
apply (clarify)
apply (frule-tac as=as in sorted-spvec-cons1)
apply (frule-tac as=bs in sorted-spvec-cons1)
apply (simp)
apply (subst sorted-spvec-step)
apply (simp split: list.split)
apply (clarify, simp)
apply (drule-tac sorted-add-spmat-helper)
apply (auto)
apply (case-tac as)
apply (simp-all)
apply (drule sorted-spvec-cons3)

```

```

apply (simp)
apply (case-tac bs)
apply (simp-all)
apply (drule sorted-spvec-cons3)
apply (simp)
done

```

```

lemma sorted-spmat-add-spmat[rule-format]: sorted-spmat A  $\longrightarrow$  sorted-spmat B
 $\longrightarrow$  sorted-spmat (add-spmat (A, B))
apply (rule add-spmat.induct[of - A B])
apply (simp-all add: sorted-spvec-add-spvec)
done

```

consts

```

le-spvec :: ('a::lordered-ab-group) spvec * 'a spvec  $\Rightarrow$  bool
le-spmat :: ('a::lordered-ab-group) spmat * 'a spmat  $\Rightarrow$  bool

```

```

recdef le-spvec measure (% (a,b). (length a) + (length b))
le-spvec ([], []) = True
le-spvec (a#as, []) = ((snd a <= 0) & (le-spvec (as, [])))
le-spvec ([], b#bs) = ((0 <= snd b) & (le-spvec ([], bs)))
le-spvec (a#as, b#bs) = (
  if (fst a < fst b) then
    ((snd a <= 0) & (le-spvec (as, b#bs)))
  else (if (fst b < fst a) then
    ((0 <= snd b) & (le-spvec (a#as, bs)))
  else
    ((snd a <= snd b) & (le-spvec (as, bs))))))

```

```

recdef le-spmat measure (% (a,b). (length a) + (length b))
le-spmat ([], []) = True
le-spmat (a#as, []) = (le-spvec (snd a, []) & (le-spmat (as, [])))
le-spmat ([], b#bs) = (le-spvec ([], snd b) & (le-spmat ([], bs)))
le-spmat (a#as, b#bs) = (
  if fst a < fst b then
    (le-spvec(snd a,[]) & le-spmat(as, b#bs))
  else (if (fst b < fst a) then
    (le-spvec([], snd b) & le-spmat(a#as, bs))
  else
    (le-spvec(snd a, snd b) & le-spmat (as, bs))))

```

constdefs

```

disj-matrices :: ('a::zero) matrix  $\Rightarrow$  'a matrix  $\Rightarrow$  bool
disj-matrices A B == (! j i. (Rep-matrix A j i  $\neq$  0)  $\longrightarrow$  (Rep-matrix B j i =
0)) & (! j i. (Rep-matrix B j i  $\neq$  0)  $\longrightarrow$  (Rep-matrix A j i = 0))

```

ML $\langle\langle$ simp-depth-limit := 6 $\rangle\rangle$

```

lemma disj-matrices-contr1: disj-matrices A B  $\Longrightarrow$  Rep-matrix A j i  $\neq$  0  $\Longrightarrow$ 

```

Rep-matrix $B j i = 0$
by (*simp add: disj-matrices-def*)

lemma *disj-matrices-contr2*: *disj-matrices* $A B \implies$ *Rep-matrix* $B j i \neq 0 \implies$
Rep-matrix $A j i = 0$
by (*simp add: disj-matrices-def*)

lemma *disj-matrices-add*: *disj-matrices* $A B \implies$ *disj-matrices* $C D \implies$ *disj-matrices*
 $A D \implies$ *disj-matrices* $B C \implies$
 $(A + B \leq C + D) = (A \leq C \ \& \ B \leq (D::('a::lordered-ab-group) \text{matrix}))$
apply (*auto*)
apply (*simp (no-asm-use) only: le-matrix-def disj-matrices-def*)
apply (*intro strip*)
apply (*erule conjE*)
apply (*drule-tac j=j and i=i in spec2*)
apply (*case-tac Rep-matrix B j i = 0*)
apply (*case-tac Rep-matrix D j i = 0*)
apply (*simp-all*)
apply (*simp (no-asm-use) only: le-matrix-def disj-matrices-def*)
apply (*intro strip*)
apply (*erule conjE*)
apply (*drule-tac j=j and i=i in spec2*)
apply (*case-tac Rep-matrix A j i = 0*)
apply (*case-tac Rep-matrix C j i = 0*)
apply (*simp-all*)
apply (*erule add-mono*)
apply (*assumption*)
done

lemma *disj-matrices-zero1*[*simp*]: *disj-matrices* $0 B$
by (*simp add: disj-matrices-def*)

lemma *disj-matrices-zero2*[*simp*]: *disj-matrices* $A 0$
by (*simp add: disj-matrices-def*)

lemma *disj-matrices-commute*: *disj-matrices* $A B =$ *disj-matrices* $B A$
by (*auto simp add: disj-matrices-def*)

lemma *disj-matrices-add-le-zero*: *disj-matrices* $A B \implies$
 $(A + B \leq 0) = (A \leq 0 \ \& \ (B::('a::lordered-ab-group) \text{matrix}) \leq 0)$
by (*rule disj-matrices-add[of A B 0 0, simplified]*)

lemma *disj-matrices-add-zero-le*: *disj-matrices* $A B \implies$
 $(0 \leq A + B) = (0 \leq A \ \& \ 0 \leq (B::('a::lordered-ab-group) \text{matrix}))$
by (*rule disj-matrices-add[of 0 0 A B, simplified]*)

lemma *disj-matrices-add-x-le*: *disj-matrices* $A B \implies$ *disj-matrices* $B C \implies$
 $(A \leq B + C) = (A \leq C \ \& \ 0 \leq (B::('a::lordered-ab-group) \text{matrix}))$

by (*auto simp add: disj-matrices-add*[of $0 A B C$, *simplified*])

lemma *disj-matrices-add-le-x*: $disj\text{-matrices } A B \implies disj\text{-matrices } B C \implies$
 $(B + A \leq C) = (A \leq C \ \& \ (B::('a::lordered\text{-ab}\text{-group})\ matrix) \leq 0)$
by (*auto simp add: disj-matrices-add*[of $B A 0 C$, *simplified*] *disj-matrices-commute*)

lemma *disj-sparse-row-singleton*: $i \leq j \implies sorted\text{-spvec}((j,y)\#v) \implies disj\text{-matrices}$
 $(sparse\text{-row}\text{-vector } v) (singleton\text{-matrix } 0 i x)$
apply (*simp add: disj-matrices-def*)
apply (*rule conjI*)
apply (*rule neg-imp*)
apply (*simp*)
apply (*intro strip*)
apply (*rule sorted-sparse-row-vector-zero*)
apply (*simp-all*)
apply (*intro strip*)
apply (*rule sorted-sparse-row-vector-zero*)
apply (*simp-all*)
done

lemma *disj-matrices-x-add*: $disj\text{-matrices } A B \implies disj\text{-matrices } A C \implies disj\text{-matrices}$
 $(A::('a::lordered\text{-ab}\text{-group})\ matrix) (B+C)$
apply (*simp add: disj-matrices-def*)
apply (*auto*)
apply (*drule-tac j=j and i=i in spec2*)+
apply (*case-tac Rep-matrix B j i = 0*)
apply (*case-tac Rep-matrix C j i = 0*)
apply (*simp-all*)
done

lemma *disj-matrices-add-x*: $disj\text{-matrices } A B \implies disj\text{-matrices } A C \implies disj\text{-matrices}$
 $(B+C) (A::('a::lordered\text{-ab}\text{-group})\ matrix)$
by (*simp add: disj-matrices-x-add disj-matrices-commute*)

lemma *disj-singleton-matrices*[*simp*]: $disj\text{-matrices } (singleton\text{-matrix } j i x) (singleton\text{-matrix}$
 $u v y) = (j \neq u \mid i \neq v \mid x = 0 \mid y = 0)$
by (*auto simp add: disj-matrices-def*)

lemma *disj-move-sparse-vec-mat*[*simplified disj-matrices-commute*]:
 $j \leq a \implies sorted\text{-spvec}((a,c)\#as) \implies disj\text{-matrices } (move\text{-matrix } (sparse\text{-row}\text{-vector}$
 $b) (int j) i) (sparse\text{-row}\text{-matrix } as)$
apply (*auto simp add: neg-def disj-matrices-def*)
apply (*drule nrows-notzero*)
apply (*drule less-le-trans*[*OF - nrows-spvec*])
apply (*subgoal-tac ja = j*)
apply (*simp add: sorted-sparse-row-matrix-zero*)
apply (*arith*)
apply (*rule nrows*)
apply (*rule order-trans*[of $- 1 -$])

```

apply (simp)
apply (case-tac nat (int ja - int j) = 0)
apply (case-tac ja = j)
apply (simp add: sorted-sparse-row-matrix-zero)
apply arith+
done

```

```

lemma disj-move-sparse-row-vector-twice:
   $j \neq u \implies \text{disj-matrices } (\text{move-matrix } (\text{sparse-row-vector } a) j i) (\text{move-matrix } (\text{sparse-row-vector } b) u v)$ 
  apply (auto simp add: neg-def disj-matrices-def)
  apply (rule nrows, rule order-trans[of - 1], simp, drule nrows-notzero, drule less-le-trans[OF - nrows-spvec], arith)+
done

```

```

lemma le-spvec-iff-sparse-row-le[rule-format]:  $(\text{sorted-spvec } a) \longrightarrow (\text{sorted-spvec } b) \longrightarrow (\text{le-spvec } (a,b)) = (\text{sparse-row-vector } a \leq \text{sparse-row-vector } b)$ 
apply (rule le-spvec.induct)
apply (simp-all add: sorted-spvec-cons1 disj-matrices-add-le-zero disj-matrices-add-zero-le

```

```

  disj-sparse-row-singleton[OF order-refl] disj-matrices-commute)
apply (rule conjI, intro strip)
apply (simp add: sorted-spvec-cons1)
apply (subst disj-matrices-add-x-le)
apply (simp add: disj-sparse-row-singleton[OF less-imp-le] disj-matrices-x-add disj-matrices-commute)
apply (simp add: disj-sparse-row-singleton[OF order-refl] disj-matrices-commute)
apply (simp, blast)
apply (intro strip, rule conjI, intro strip)
apply (simp add: sorted-spvec-cons1)
apply (subst disj-matrices-add-le-x)
apply (simp-all add: disj-sparse-row-singleton[OF order-refl] disj-sparse-row-singleton[OF less-imp-le] disj-matrices-commute disj-matrices-x-add)
apply (blast)
apply (intro strip)
apply (simp add: sorted-spvec-cons1)
apply (case-tac a=aa, simp-all)
apply (subst disj-matrices-add)
apply (simp-all add: disj-sparse-row-singleton[OF order-refl] disj-matrices-commute)
done

```

```

lemma le-spvec-empty2-sparse-row[rule-format]:  $(\text{sorted-spvec } b) \longrightarrow (\text{le-spvec } (b, [])) = (\text{sparse-row-vector } b \leq 0)$ 
apply (induct b)
apply (simp-all add: sorted-spvec-cons1)
apply (intro strip)
apply (subst disj-matrices-add-le-zero)
apply (simp add: disj-matrices-commute disj-sparse-row-singleton sorted-spvec-cons1)
apply (rule-tac y = snd a in disj-sparse-row-singleton[OF order-refl])

```

apply (*simp-all*)
done

lemma *le-spvec-empty1-sparse-row*[*rule-format*]: (*sorted-spvec b*) \longrightarrow (*le-spvec* (\square , *b*)
 $=$ ($0 \leq$ *sparse-row-vector b*))
apply (*induct b*)
apply (*simp-all add: sorted-spvec-cons1*)
apply (*intro strip*)
apply (*subst disj-matrices-add-zero-le*)
apply (*simp add: disj-matrices-commute disj-sparse-row-singleton sorted-spvec-cons1*)
apply (*rule-tac y = snd a in disj-sparse-row-singleton[OF order-refl]*)
apply (*simp-all*)
done

lemma *le-spmat-iff-sparse-row-le*[*rule-format*]: (*sorted-spvec A*) \longrightarrow (*sorted-spmat*
A) \longrightarrow (*sorted-spvec B*) \longrightarrow (*sorted-spmat B*) \longrightarrow
le-spmat(A, B) = (sparse-row-matrix A \leq *sparse-row-matrix B)*
apply (*rule le-spmat.induct*)
apply (*simp add: sparse-row-matrix-cons disj-matrices-add-le-zero disj-matrices-add-zero-le*
disj-move-sparse-vec-mat[OF order-refl]
disj-matrices-commute sorted-spvec-cons1 le-spvec-empty2-sparse-row le-spvec-empty1-sparse-row)
done

apply (*rule conjI, intro strip*)
apply (*simp add: sorted-spvec-cons1*)
apply (*subst disj-matrices-add-x-le*)
apply (*rule disj-matrices-add-x*)
apply (*simp add: disj-move-sparse-row-vector-twice*)
apply (*simp add: disj-move-sparse-vec-mat[OF less-imp-le] disj-matrices-commute*)
apply (*simp add: disj-move-sparse-vec-mat[OF order-refl] disj-matrices-commute*)
apply (*simp, blast*)
apply (*intro strip, rule conjI, intro strip*)
apply (*simp add: sorted-spvec-cons1*)
apply (*subst disj-matrices-add-le-x*)
apply (*simp add: disj-move-sparse-vec-mat[OF order-refl]*)
apply (*rule disj-matrices-x-add*)
apply (*simp add: disj-move-sparse-row-vector-twice*)
apply (*simp add: disj-move-sparse-vec-mat[OF less-imp-le] disj-matrices-commute*)
apply (*simp, blast*)
apply (*intro strip*)
apply (*case-tac a=aa*)
apply (*simp-all*)
apply (*subst disj-matrices-add*)
apply (*simp-all add: disj-matrices-commute disj-move-sparse-vec-mat[OF order-refl]*)
apply (*simp add: sorted-spvec-cons1 le-spvec-iff-sparse-row-le*)
done

ML $\langle\langle$ *simp-depth-limit := 999* $\rangle\rangle$

consts

$abs\text{-spmat} :: ('a::\text{lordered-ring})\ spmat \Rightarrow 'a\ spmat$
 $minus\text{-spmat} :: ('a::\text{lordered-ring})\ spmat \Rightarrow 'a\ spmat$

primrec

$abs\text{-spmat}\ [] = []$
 $abs\text{-spmat}\ (a\#as) = (fst\ a,\ abs\text{-spvec}\ (snd\ a))\#(abs\text{-spmat}\ as)$

primrec

$minus\text{-spmat}\ [] = []$
 $minus\text{-spmat}\ (a\#as) = (fst\ a,\ minus\text{-spvec}\ (snd\ a))\#(minus\text{-spmat}\ as)$

lemma *sparse-row-matrix-minus:*

$sparse\text{-row-matrix}\ (minus\text{-spmat}\ A) = -\ (sparse\text{-row-matrix}\ A)$
apply (*induct* A)
apply (*simp-all add: sparse-row-vector-minus sparse-row-matrix-cons*)
apply (*subst Rep-matrix-inject[symmetric]*)
apply (*rule ext*)
apply *simp*
done

lemma *Rep-sparse-row-vector-zero: $x \neq 0 \implies Rep\text{-matrix}\ (sparse\text{-row-vector}\ v)$*
 $x\ y = 0$

proof –

assume $x:x \neq 0$
have $r:nrows\ (sparse\text{-row-vector}\ v) \leq Suc\ 0$ **by** (*rule nrows-spvec*)
show *?thesis*
apply (*rule nrows*)
apply (*subgoal-tac Suc 0 <= x*)
apply (*insert r*)
apply (*simp only:*)
apply (*insert x*)
apply *arith*
done

qed

lemma *sparse-row-matrix-abs:*

$sorted\text{-spvec}\ A \implies sorted\text{-spmat}\ A \implies sparse\text{-row-matrix}\ (abs\text{-spmat}\ A) = abs$
 $(sparse\text{-row-matrix}\ A)$
apply (*induct* A)
apply (*simp-all add: sparse-row-vector-abs sparse-row-matrix-cons*)
apply (*frule-tac sorted-spvec-cons1, simp*)
apply (*simplesubst Rep-matrix-inject[symmetric]*)
apply (*rule ext*)
apply *auto*
apply (*case-tac x=a*)
apply (*simp*)
apply (*simplesubst sorted-sparse-row-matrix-zero*)
apply *auto*
apply (*simplesubst Rep-sparse-row-vector-zero*)

```

apply (simp-all add: neg-def)
done

lemma sorted-spvec-minus-spmat: sorted-spvec A  $\implies$  sorted-spvec (minus-spmat A)
apply (induct A)
apply (simp)
apply (frule sorted-spvec-cons1, simp)
apply (simp add: sorted-spvec.simps split:list.split-asm)
done

lemma sorted-spvec-abs-spmat: sorted-spvec A  $\implies$  sorted-spvec (abs-spmat A)
apply (induct A)
apply (simp)
apply (frule sorted-spvec-cons1, simp)
apply (simp add: sorted-spvec.simps split:list.split-asm)
done

lemma sorted-spmat-minus-spmat: sorted-spmat A  $\implies$  sorted-spmat (minus-spmat A)
apply (induct A)
apply (simp-all add: sorted-spvec-minus-spvec)
done

lemma sorted-spmat-abs-spmat: sorted-spmat A  $\implies$  sorted-spmat (abs-spmat A)
apply (induct A)
apply (simp-all add: sorted-spvec-abs-spvec)
done

constdefs
  diff-spmat :: ('a::lordered-ring) spmat  $\Rightarrow$  'a spmat  $\Rightarrow$  'a spmat
  diff-spmat A B == add-spmat (A, minus-spmat B)

lemma sorted-spmat-diff-spmat: sorted-spmat A  $\implies$  sorted-spmat B  $\implies$  sorted-spmat (diff-spmat A B)
by (simp add: diff-spmat-def sorted-spmat-minus-spmat sorted-spmat-add-spmat)

lemma sorted-spvec-diff-spmat: sorted-spvec A  $\implies$  sorted-spvec B  $\implies$  sorted-spvec (diff-spmat A B)
by (simp add: diff-spmat-def sorted-spvec-minus-spmat sorted-spvec-add-spmat)

lemma sparse-row-diff-spmat: sparse-row-matrix (diff-spmat A B) = (sparse-row-matrix A) - (sparse-row-matrix B)
by (simp add: diff-spmat-def sparse-row-add-spmat sparse-row-matrix-minus)

constdefs
  sorted-sparse-matrix :: 'a spmat  $\Rightarrow$  bool
  sorted-sparse-matrix A == (sorted-spvec A) & (sorted-spmat A)

```

lemma *sorted-sparse-matrix-imp-spvec*: *sorted-sparse-matrix* $A \implies$ *sorted-spvec* A
by (*simp add: sorted-sparse-matrix-def*)

lemma *sorted-sparse-matrix-imp-spmat*: *sorted-sparse-matrix* $A \implies$ *sorted-spmat* A
by (*simp add: sorted-sparse-matrix-def*)

lemmas *sorted-sp-simps* =
sorted-spvec.simps
sorted-spmat.simps
sorted-sparse-matrix-def

lemma *bool1*: $(\neg \text{True}) = \text{False}$ **by** *blast*

lemma *bool2*: $(\neg \text{False}) = \text{True}$ **by** *blast*

lemma *bool3*: $((P::\text{bool}) \wedge \text{True}) = P$ **by** *blast*

lemma *bool4*: $(\text{True} \wedge (P::\text{bool})) = P$ **by** *blast*

lemma *bool5*: $((P::\text{bool}) \wedge \text{False}) = \text{False}$ **by** *blast*

lemma *bool6*: $(\text{False} \wedge (P::\text{bool})) = \text{False}$ **by** *blast*

lemma *bool7*: $((P::\text{bool}) \vee \text{True}) = \text{True}$ **by** *blast*

lemma *bool8*: $(\text{True} \vee (P::\text{bool})) = \text{True}$ **by** *blast*

lemma *bool9*: $((P::\text{bool}) \vee \text{False}) = P$ **by** *blast*

lemma *bool10*: $(\text{False} \vee (P::\text{bool})) = P$ **by** *blast*

lemmas *boolarith* = *bool1 bool2 bool3 bool4 bool5 bool6 bool7 bool8 bool9 bool10*

lemma *if-case-eq*: $(\text{if } b \text{ then } x \text{ else } y) = (\text{case } b \text{ of } \text{True} \implies x \mid \text{False} \implies y)$ **by** *simp*

consts

pprt-spvec :: $('a::\{\text{lordered-ab-group}\}) \text{ spvec} \Rightarrow 'a \text{ spvec}$

nprrt-spvec :: $('a::\{\text{lordered-ab-group}\}) \text{ spvec} \Rightarrow 'a \text{ spvec}$

pprt-spmat :: $('a::\{\text{lordered-ab-group}\}) \text{ spmat} \Rightarrow 'a \text{ spmat}$

nprrt-spmat :: $('a::\{\text{lordered-ab-group}\}) \text{ spmat} \Rightarrow 'a \text{ spmat}$

primrec

pprt-spvec $[] = []$

pprt-spvec $(a\#as) = (\text{fst } a, \text{pprt } (\text{snd } a)) \# (\text{pprt-spvec } as)$

primrec

nprrt-spvec $[] = []$

nprrt-spvec $(a\#as) = (\text{fst } a, \text{nprrt } (\text{snd } a)) \# (\text{nprrt-spvec } as)$

primrec

pprt-spmat $[] = []$

pprt-spmat $(a\#as) = (\text{fst } a, \text{pprt-spvec } (\text{snd } a)) \# (\text{pprt-spmat } as)$

primrec

nprrt-spmat $[] = []$

nprrt-spmat $(a\#as) = (\text{fst } a, \text{nprrt-spvec } (\text{snd } a)) \# (\text{nprrt-spmat } as)$

lemma *pprt-add: disj-matrices* A ($B::(\text{::}:\text{ordered-ring})$ *matrix*) \implies *pprt* ($A+B$)
 $=$ *pprt* A + *pprt* B
apply (*simp add: pprt-def join-matrix*)
apply (*simp add: Rep-matrix-inject[symmetric]*)
apply (*rule ext*) +
apply *simp*
apply (*case-tac Rep-matrix A x xa $\neq 0$*)
apply (*simp-all add: disj-matrices-contr1*)
done

lemma *nprt-add: disj-matrices* A ($B::(\text{::}:\text{ordered-ring})$ *matrix*) \implies *nprt* ($A+B$)
 $=$ *nprt* A + *nprt* B
apply (*simp add: nprt-def meet-matrix*)
apply (*simp add: Rep-matrix-inject[symmetric]*)
apply (*rule ext*) +
apply *simp*
apply (*case-tac Rep-matrix A x xa $\neq 0$*)
apply (*simp-all add: disj-matrices-contr1*)
done

lemma *pprt-singleton[simp]: pprt* (*singleton-matrix* j i ($x::(\text{::}:\text{ordered-ring})$)) $=$ *singleton-matrix* j i (*pprt* x)
apply (*simp add: pprt-def join-matrix*)
apply (*simp add: Rep-matrix-inject[symmetric]*)
apply (*rule ext*) +
apply *simp*
done

lemma *nprt-singleton[simp]: nprt* (*singleton-matrix* j i ($x::(\text{::}:\text{ordered-ring})$)) $=$ *singleton-matrix* j i (*nprt* x)
apply (*simp add: nprt-def meet-matrix*)
apply (*simp add: Rep-matrix-inject[symmetric]*)
apply (*rule ext*) +
apply *simp*
done

lemma *less-imp-le: a < b \implies a <= (b::(\text{::}:\text{order}))* **by** (*simp add: less-def*)

lemma *sparse-row-vector-pprt: sorted-spvec* v \implies *sparse-row-vector* (*pprt-spvec* v) $=$ *pprt* (*sparse-row-vector* v)
apply (*induct v*)
apply (*simp-all*)
apply (*frule sorted-spvec-cons1, auto*)
apply (*subst pprt-add*)
apply (*subst disj-matrices-commute*)
apply (*rule disj-sparse-row-singleton*)

apply *auto*
done

lemma *sparse-row-vector-nprt*: $\text{sorted-spvec } v \implies \text{sparse-row-vector } (\text{nprt-spvec } v) = \text{nprt } (\text{sparse-row-vector } v)$
apply (*induct* *v*)
apply (*simp-all*)
apply (*frule sorted-spvec-cons1*, *auto*)
apply (*subst nprt-add*)
apply (*subst disj-matrices-commute*)
apply (*rule disj-sparse-row-singleton*)
apply *auto*
done

lemma *pprt-move-matrix*: $\text{pprt } (\text{move-matrix } (A::('a::\text{lordered-ring}) \text{matrix}) \text{ } j \text{ } i) = \text{move-matrix } (\text{pprt } A) \text{ } j \text{ } i$
apply (*simp add: pprt-def*)
apply (*simp add: join-matrix*)
apply (*simp add: Rep-matrix-inject[symmetric]*)
apply (*rule ext*)+
apply (*simp*)
done

lemma *nprt-move-matrix*: $\text{nprt } (\text{move-matrix } (A::('a::\text{lordered-ring}) \text{matrix}) \text{ } j \text{ } i) = \text{move-matrix } (\text{nprt } A) \text{ } j \text{ } i$
apply (*simp add: nprt-def*)
apply (*simp add: meet-matrix*)
apply (*simp add: Rep-matrix-inject[symmetric]*)
apply (*rule ext*)+
apply (*simp*)
done

lemma *sparse-row-matrix-pprt*: $\text{sorted-spvec } m \implies \text{sorted-spmat } m \implies \text{sparse-row-matrix } (\text{pprt-spmat } m) = \text{pprt } (\text{sparse-row-matrix } m)$
apply (*induct* *m*)
apply *simp*
apply *simp*
apply (*frule sorted-spvec-cons1*)
apply (*simp add: sparse-row-matrix-cons sparse-row-vector-pprt*)
apply (*subst pprt-add*)
apply (*subst disj-matrices-commute*)
apply (*rule disj-move-sparse-vec-mat*)
apply *auto*
apply (*simp add: sorted-spvec.simps*)
apply (*simp split: list.split*)
apply *auto*
apply (*simp add: pprt-move-matrix*)
done

lemma *sparse-row-matrix-nprt: sorted-spvec m \implies sorted-spmat m \implies sparse-row-matrix (nprt-spmat m) = nprt (sparse-row-matrix m)*

```

apply (induct m)
apply simp
apply simp
apply (frule sorted-spvec-cons1)
apply (simp add: sparse-row-matrix-cons sparse-row-vector-nprt)
apply (subst nprt-add)
apply (subst disj-matrices-commute)
apply (rule disj-move-sparse-vec-mat)
apply auto
apply (simp add: sorted-spvec.simps)
apply (simp split: list.split)
apply auto
apply (simp add: nprt-move-matrix)
done

```

lemma *sorted-pprt-spvec: sorted-spvec v \implies sorted-spvec (pprt-spvec v)*

```

apply (induct v)
apply (simp)
apply (frule sorted-spvec-cons1)
apply simp
apply (simp add: sorted-spvec.simps split:list.split-asm)
done

```

lemma *sorted-nprt-spvec: sorted-spvec v \implies sorted-spvec (nprt-spvec v)*

```

apply (induct v)
apply (simp)
apply (frule sorted-spvec-cons1)
apply simp
apply (simp add: sorted-spvec.simps split:list.split-asm)
done

```

lemma *sorted-spvec-pprt-spmat: sorted-spvec m \implies sorted-spvec (pprt-spmat m)*

```

apply (induct m)
apply (simp)
apply (frule sorted-spvec-cons1)
apply simp
apply (simp add: sorted-spvec.simps split:list.split-asm)
done

```

lemma *sorted-spvec-nprt-spmat: sorted-spvec m \implies sorted-spvec (nprt-spmat m)*

```

apply (induct m)
apply (simp)
apply (frule sorted-spvec-cons1)
apply simp
apply (simp add: sorted-spvec.simps split:list.split-asm)
done

```

lemma *sorted-spmat-pprt-spmat*: *sorted-spmat m* \implies *sorted-spmat (pprt-spmat m)*
apply (*induct m*)
apply (*simp-all add: sorted-pprt-spmat*)
done

lemma *sorted-spmat-nprt-spmat*: *sorted-spmat m* \implies *sorted-spmat (nprt-spmat m)*
apply (*induct m*)
apply (*simp-all add: sorted-nprt-spmat*)
done

constdefs

mult-est-spmat :: ('a::lordered-ring) *spmat* \Rightarrow 'a *spmat* \Rightarrow 'a *spmat* \Rightarrow 'a *spmat*
 \Rightarrow 'a *spmat*
mult-est-spmat r1 r2 s1 s2 ==
add-spmat (mult-spmat (pprt-spmat s2) (pprt-spmat r2), add-spmat (mult-spmat (pprt-spmat s1) (nprt-spmat r2), add-spmat (mult-spmat (nprt-spmat s2) (pprt-spmat r1), mult-spmat (nprt-spmat s1) (nprt-spmat r1))))

lemmas *sparse-row-matrix-op-simps* =

sorted-sparse-matrix-imp-spmat sorted-sparse-matrix-imp-spmat
sparse-row-add-spmat sorted-spmat-add-spmat sorted-spmat-add-spmat
sparse-row-diff-spmat sorted-spmat-diff-spmat sorted-spmat-diff-spmat
sparse-row-matrix-minus sorted-spmat-minus-spmat sorted-spmat-minus-spmat
sparse-row-mult-spmat sorted-spmat-mult-spmat sorted-spmat-mult-spmat
sparse-row-matrix-abs sorted-spmat-abs-spmat sorted-spmat-abs-spmat
le-spmat-iff-sparse-row-le
sparse-row-matrix-pprt sorted-spmat-pprt-spmat sorted-spmat-pprt-spmat
sparse-row-matrix-nprt sorted-spmat-nprt-spmat sorted-spmat-nprt-spmat

lemma *zero-eq-Numeral0*: (0:::number-ring) = *Numeral0* **by** *simp*

lemmas *sparse-row-matrix-arith-simps*[*simplified zero-eq-Numeral0*] =

mult-spmat.simps mult-spmat.simps
addmult-spmat.simps
smult-spmat-empty smult-spmat-cons
add-spmat.simps add-spmat.simps
minus-spmat.simps minus-spmat.simps
abs-spmat.simps abs-spmat.simps
diff-spmat-def
le-spmat.simps le-spmat.simps
pprt-spmat.simps pprt-spmat.simps
nprt-spmat.simps nprt-spmat.simps
mult-est-spmat-def

lemma *spm-mult-le-dual-prts*:

assumes

sorted-sparse-matrix $A1$

sorted-sparse-matrix $A2$

sorted-sparse-matrix $c1$

sorted-sparse-matrix $c2$

sorted-sparse-matrix y

sorted-sparse-matrix $r1$

sorted-sparse-matrix $r2$

sorted-spvec b

le-spmat (\square , y)

sparse-row-matrix $A1 \leq A$

$A \leq$ *sparse-row-matrix* $A2$

sparse-row-matrix $c1 \leq c$

$c \leq$ *sparse-row-matrix* $c2$

sparse-row-matrix $r1 \leq x$

$x \leq$ *sparse-row-matrix* $r2$

$A * x \leq$ *sparse-row-matrix* ($b::('a::lordered-ring)$ *spmat*)

shows

$c * x \leq$ *sparse-row-matrix* (*add-spmat* (*mult-spmat* y b ,

(*let* $s1 =$ *diff-spmat* $c1$ (*mult-spmat* y $A2$); $s2 =$ *diff-spmat* $c2$ (*mult-spmat* y $A1$) *in*

add-spmat (*mult-spmat* (*pprt-spmat* $s2$) (*pprt-spmat* $r2$), *add-spmat* (*mult-spmat* (*pprt-spmat* $s1$) (*nprrt-spmat* $r2$),

add-spmat (*mult-spmat* (*nprrt-spmat* $s2$) (*pprt-spmat* $r1$), *mult-spmat* (*nprrt-spmat* $s1$) (*nprrt-spmat* $r1$))))))

apply (*simp* *add: Let-def*)

apply (*insert prems*)

apply (*simp* *add: sparse-row-matrix-op-simps ring-eq-simps*)

apply (*rule* *mult-le-dual-prts*[**where** $A=A$, *simplified Let-def ring-eq-simps*])

apply (*auto*)

done

lemma *spm-mult-le-dual-prts-no-let*:

assumes

sorted-sparse-matrix $A1$

sorted-sparse-matrix $A2$

sorted-sparse-matrix $c1$

sorted-sparse-matrix $c2$

sorted-sparse-matrix y

sorted-sparse-matrix $r1$

sorted-sparse-matrix $r2$

sorted-spvec b

le-spmat (\square , y)

sparse-row-matrix $A1 \leq A$

$A \leq$ *sparse-row-matrix* $A2$

sparse-row-matrix $c1 \leq c$

```

  c ≤ sparse-row-matrix c2
  sparse-row-matrix r1 ≤ x
  x ≤ sparse-row-matrix r2
  A * x ≤ sparse-row-matrix (b::('a::lordered-ring) spmat)
shows
  c * x ≤ sparse-row-matrix (add-spmat (mult-spmat y b,
  mult-est-spmat r1 r2 (diff-spmat c1 (mult-spmat y A2)) (diff-spmat c2 (mult-spmat
  y A1))))
  by (simp add: prems mult-est-spmat-def spm-mult-le-dual-prts[where A=A, sim-
  plified Let-def])

end

```

```

theory FloatSparseMatrix imports Float SparseMatrix begin

end

```

```

theory Cplex
imports FloatSparseMatrix
uses Cplex-tools.ML CplexMatrixConverter.ML FloatSparseMatrixBuilder.ML fspmlp.ML
begin

end

```

```

theory MatrixLP
imports Cplex
begin

```

```

constdefs
  list-case-compute :: 'b list ⇒ 'a ⇒ ('b ⇒ 'b list ⇒ 'a) ⇒ 'a
  list-case-compute l a f == list-case a f l

```

```

lemma list-case-compute: list-case = (λ (a::'a) f (l::'b list). list-case-compute l a
f)
  apply (rule ext)+
  apply (simp add: list-case-compute-def)
  done

```

```

lemma list-case-compute-empty: list-case-compute ([]::'b list) = (λ (a::'a) f. a)
  apply (rule ext)+

```

```

apply (simp add: list-case-compute-def)
done

lemma list-case-compute-cons: list-case-compute (u#v) = (λ (a::'a) f. (f (u::'b)
v))
apply (rule ext)+
apply (simp add: list-case-compute-def)
done

lemma If-True: (If True) = (λ x y. x)
apply (rule ext)+
apply auto
done

lemma If-False: (If False) = (λ x y. y)
apply (rule ext)+
apply auto
done

lemma Let-compute: Let (x::'a) f = ((f x)::'b)
by (simp add: Let-def)

lemma fst-compute: fst (a::'a, b::'b) = a
by auto

lemma snd-compute: snd (a::'a, b::'b) = b
by auto

lemma bool1: (¬ True) = False by blast
lemma bool2: (¬ False) = True by blast
lemma bool3: ((P::bool) ∧ True) = P by blast
lemma bool4: (True ∧ (P::bool)) = P by blast
lemma bool5: ((P::bool) ∧ False) = False by blast
lemma bool6: (False ∧ (P::bool)) = False by blast
lemma bool7: ((P::bool) ∨ True) = True by blast
lemma bool8: (True ∨ (P::bool)) = True by blast
lemma bool9: ((P::bool) ∨ False) = P by blast
lemma bool10: (False ∨ (P::bool)) = P by blast
lemmas boolarith = bool1 bool2 bool3 bool4 bool5 bool6 bool7 bool8 bool9 bool10

lemmas float-arith = Float.arith
lemmas sparse-row-matrix-arith-simps = SparseMatrix.sparse-row-matrix-arith-simps
lemmas sorted-sp-simps = SparseMatrix.sorted-sp-simps
lemmas fst-snd-conv = Product-Type.fst-conv Product-Type.snd-conv

end

```