

IMP in HOLCF

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1 Denotational Semantics of Commands in HOLCF

theory *Denotational* imports *HOLCF Natural* begin

1.1 Definition

constdefs

```
dlift :: "('a::type) discr -> 'b::pcpo) => ('a lift -> 'b)"  
"dlift f == (LAM x. case x of UU => UU | Def y => f.(Discr y))"
```

```
consts D :: "com => state discr -> state lift"
```

primrec

```
"D(skip) = (LAM s. Def(undiscr s))"  
"D(X ::= a) = (LAM s. Def((undiscr s)[X ↦ a(undiscr s)]))"  
"D(c0 ; c1) = (dlift(D c1) oo (D c0))"  
"D(if b then c1 else c2) =  
  (LAM s. if b (undiscr s) then (D c1)·s else (D c2)·s)"  
"D(while b do c) =  
  fix·(LAM w s. if b (undiscr s) then (dlift w)·((D c)·s)  
    else Def(undiscr s))"
```

1.2 Equivalence of Denotational Semantics in HOLCF and Evaluation Semantics in HOL

```
lemma dlift_Def [simp]: "dlift f·(Def x) = f·(Discr x)"  
  <proof>
```

```

lemma cont_dlift [iff]: "cont (%f. dlift f)"
  <proof>

lemma dlift_is_Def [simp]:
  "(dlift f.l = Def y) = ( $\exists x. l = Def x \wedge f.(Discr x) = Def y$ )"
  <proof>

lemma eval_implies_D: "<math>\langle c, s \rangle \longrightarrow_c t \implies D c.(Discr s) = (Def t)</math>"
  <proof>

lemma D_implies_eval: "!s t. D c.(Discr s) = (Def t) --> <math>\langle c, s \rangle \longrightarrow_c t</math>"
  <proof>

theorem D_is_eval: "(D c.(Discr s) = (Def t)) = (<math>\langle c, s \rangle \longrightarrow_c t</math>)"
  <proof>

end

```

2 Correctness of Hoare by Fixpoint Reasoning

```

theory HoareEx imports Denotational begin

```

An example from the HOLCF paper by Müller, Nipkow, Oheimb, Slotosch [1]. It demonstrates fixpoint reasoning by showing the correctness of the Hoare rule for while-loops.

```

types assn = "state => bool"

```

```

constdefs

```

```

  hoare_valid :: "[assn, com, assn] => bool"    ("|= {(1_)} / (_)/ {(1_)}" 50)
  "|= {A} c {B} ==  $\forall s t. A s \wedge D c \$ (Discr s) = Def t \longrightarrow B t$ "

```

```

lemma WHILE_rule_sound:

```

```

  "|= {A} c {A} ==> |= {A} while b do c { $\lambda s. A s \wedge \neg b s$ }"
  <proof>

```

```

end

```

References

- [1] O. Müller, T. Nipkow, D. v. Oheimb, and O. Slotosch. HOLCF = HOL + LCF. *J. Functional Programming*, 9:191–223, 1999.