

IMP in HOLCF

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1 Denotational Semantics of Commands in HOLCF

theory *Denotational* imports *HOLCF* *Natural* begin

1.1 Definition

constdefs

```
dlift :: "('a::type) discr -> 'b::pcpo) => ('a lift -> 'b)"
"dlift f == (LAM x. case x of UU => UU | Def y => f.(Discr y))"
```

consts *D* :: "com => state discr -> state lift"

primrec

```
"D(skip) = (LAM s. Def(undiscr s))"
"D(X ::= a) = (LAM s. Def((undiscr s)[X ↦ a(undiscr s)]))"
"D(c0 ; c1) = (dlift(D c1) oo (D c0))"
"D(if b then c1 else c2) =
  (LAM s. if b (undiscr s) then (D c1)·s else (D c2)·s)"
"D(while b do c) =
  fix·(LAM w s. if b (undiscr s) then (dlift w)·((D c)·s)
    else Def(undiscr s))"
```

1.2 Equivalence of Denotational Semantics in HOLCF and Evaluation Semantics in HOL

lemma *dlift_Def [simp]*: "dlift f.(Def x) = f.(Discr x)"
⟨*proof*⟩

```

lemma cont_dlift [iff]: "cont (%f. dlift f)"
  <proof>

lemma dlift_is_Def [simp]:
  "(dlift f.l = Def y) = ( $\exists x. l = \text{Def } x \wedge f \cdot (\text{Discr } x) = \text{Def } y$ )"
  <proof>

lemma eval_implies_D: " $\langle c, s \rangle \longrightarrow_c t \implies D \ c \cdot (\text{Discr } s) = (\text{Def } t)$ "
  <proof>

lemma D_implies_eval: " $!s \ t. D \ c \cdot (\text{Discr } s) = (\text{Def } t) \implies \langle c, s \rangle \longrightarrow_c t$ "
  <proof>

theorem D_is_eval: " $(D \ c \cdot (\text{Discr } s) = (\text{Def } t)) = (\langle c, s \rangle \longrightarrow_c t)$ "
  <proof>

end

```

2 Correctness of Hoare by Fixpoint Reasoning

theory HoareEx imports Denotational begin

An example from the HOLCF paper by Müller, Nipkow, Oheimb, Slotosch [1]. It demonstrates fixpoint reasoning by showing the correctness of the Hoare rule for while-loops.

```

types assn = "state => bool"

constdefs
  hoare_valid :: "[assn, com, assn] => bool"    ("|= {(1_)} / (_)/ {(1_)}" 50)
  "|= {A} c {B} ==  $\forall s \ t. A \ s \wedge D \ c \ \$ (\text{Discr } s) = \text{Def } t \implies B \ t$ "

lemma WHILE_rule_sound:
  " $|= \{A\} \ c \ \{A\} \implies \text{|= } \{A\} \ \text{while } b \ \text{do } c \ \{\lambda s. A \ s \wedge \neg b \ s\}$ "
  <proof>

end

```

References

- [1] O. Müller, T. Nipkow, D. v. Oheimb, and O. Slotosch. HOLCF = HOL + LCF. *J. Functional Programming*, 9:191–223, 1999.