

# Matrix

Steven Obua

October 1, 2005

```
theory MatrixGeneral imports Main begin

types 'a infmatrix = [nat, nat]  $\Rightarrow$  'a

constdefs
  nonzero-positions :: ('a::zero) infmatrix  $\Rightarrow$  (nat*nat) set
  nonzero-positions A == {pos. A (fst pos) (snd pos)  $\sim$  0}

typedef 'a matrix = {(f::('a::zero) infmatrix)}. finite (nonzero-positions f)}
<proof>

declare Rep-matrix-inverse[simp]

lemma finite-nonzero-positions : finite (nonzero-positions (Rep-matrix A))
<proof>

constdefs
  nrows :: ('a::zero) matrix  $\Rightarrow$  nat
  nrows A == if nonzero-positions(Rep-matrix A) = {} then 0 else Suc(Max
((image fst) (nonzero-positions (Rep-matrix A))))
  ncols :: ('a::zero) matrix  $\Rightarrow$  nat
  ncols A == if nonzero-positions(Rep-matrix A) = {} then 0 else Suc(Max ((image
snd) (nonzero-positions (Rep-matrix A))))

lemma nrows:
  assumes hyp: nrows A  $\leq$  m
  shows (Rep-matrix A m n) = 0 (is ?concl)
<proof>

constdefs
  transpose-infmatrix :: 'a infmatrix  $\Rightarrow$  'a infmatrix
  transpose-infmatrix A j i == A i j
  transpose-matrix :: ('a::zero) matrix  $\Rightarrow$  'a matrix
  transpose-matrix == Abs-matrix o transpose-infmatrix o Rep-matrix

declare transpose-infmatrix-def[simp]
```

**lemma** *transpose-infmatrix-twice*[simp]: *transpose-infmatrix (transpose-infmatrix A) = A*  
 ⟨proof⟩

**lemma** *transpose-infmatrix*: *transpose-infmatrix (% j i. P j i) = (% j i. P i j)*  
 ⟨proof⟩

**lemma** *transpose-infmatrix-closed*[simp]: *Rep-matrix (Abs-matrix (transpose-infmatrix (Rep-matrix x))) = transpose-infmatrix (Rep-matrix x)*  
 ⟨proof⟩

**lemma** *infmatrixforward*: *(x::'a infmatrix) = y  $\implies$   $\forall$  a b. x a b = y a b* ⟨proof⟩

**lemma** *transpose-infmatrix-inject*: *(transpose-infmatrix A = transpose-infmatrix B) = (A = B)*  
 ⟨proof⟩

**lemma** *transpose-matrix-inject*: *(transpose-matrix A = transpose-matrix B) = (A = B)*  
 ⟨proof⟩

**lemma** *transpose-matrix*[simp]: *Rep-matrix(transpose-matrix A) j i = Rep-matrix A i j*  
 ⟨proof⟩

**lemma** *transpose-transpose-id*[simp]: *transpose-matrix (transpose-matrix A) = A*  
 ⟨proof⟩

**lemma** *nrows-transpose*[simp]: *nrows (transpose-matrix A) = ncols A*  
 ⟨proof⟩

**lemma** *ncols-transpose*[simp]: *ncols (transpose-matrix A) = nrows A*  
 ⟨proof⟩

**lemma** *ncols*: *ncols A <= n  $\implies$  Rep-matrix A m n = 0*  
 ⟨proof⟩

**lemma** *ncols-le*: *(ncols A <= n) = (! j i. n <= i  $\longrightarrow$  (Rep-matrix A j i) = 0) (is - = ?st)*  
 ⟨proof⟩

**lemma** *less-ncols*: *(n < ncols A) = (? j i. n <= i & (Rep-matrix A j i)  $\neq$  0) (is ?concl)*  
 ⟨proof⟩

**lemma** *le-ncols*: *(n <= ncols A) = ( $\forall$  m. ( $\forall$  j i. m <= i  $\longrightarrow$  (Rep-matrix A j i) = 0)  $\longrightarrow$  n <= m) (is ?concl)*  
 ⟨proof⟩

**lemma** *nrows-le*:  $(nrows\ A \leq n) = (!\ j\ i.\ n \leq j \longrightarrow (Rep\ matrix\ A\ j\ i) = 0)$   
**(is ?s)**  
 <proof>

**lemma** *less-nrows*:  $(m < nrows\ A) = (?\ j\ i.\ m \leq j \ \&\ (Rep\ matrix\ A\ j\ i) \neq 0)$   
**(is ?concl)**  
 <proof>

**lemma** *le-nrows*:  $(n \leq nrows\ A) = (\forall\ m.\ (\forall\ j\ i.\ m \leq j \longrightarrow (Rep\ matrix\ A\ j\ i) = 0) \longrightarrow n \leq m)$  **(is ?concl)**  
 <proof>

**lemma** *nrows-notzero*:  $Rep\ matrix\ A\ m\ n \neq 0 \implies m < nrows\ A$   
 <proof>

**lemma** *ncols-notzero*:  $Rep\ matrix\ A\ m\ n \neq 0 \implies n < ncols\ A$   
 <proof>

**lemma** *finite-natarray1*:  $finite\ \{x.\ x < (n::nat)\}$   
 <proof>

**lemma** *finite-natarray2*:  $finite\ \{pos.\ (fst\ pos) < (m::nat) \ \&\ (snd\ pos) < (n::nat)\}$   
 <proof>

**lemma** *RepAbs-matrix*:  
**assumes** *aem*:  $?\ m.\ !\ j\ i.\ m \leq j \longrightarrow x\ j\ i = 0$  **(is ?em)** **and** *aen*:  $?\ n.\ !\ j\ i.\ (n \leq i \longrightarrow x\ j\ i = 0)$  **(is ?en)**  
**shows**  $(Rep\ matrix\ (Abs\ matrix\ x)) = x$   
 <proof>

### constdefs

*apply-infmatrix* ::  $('a \Rightarrow 'b) \Rightarrow 'a\ infmatrix \Rightarrow 'b\ infmatrix$   
*apply-infmatrix* *f* == % *A*. (% *j i*. *f* (*A j i*))  
*apply-matrix* ::  $('a \Rightarrow 'b) \Rightarrow ('a::zero)\ matrix \Rightarrow ('b::zero)\ matrix$   
*apply-matrix* *f* == % *A*. *Abs-matrix* (*apply-infmatrix* *f* (*Rep-matrix* *A*))  
*combine-infmatrix* ::  $('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow 'a\ infmatrix \Rightarrow 'b\ infmatrix \Rightarrow 'c\ infmatrix$   
*combine-infmatrix* *f* == % *A B*. (% *j i*. *f* (*A j i*) (*B j i*))  
*combine-matrix* ::  $('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow ('a::zero)\ matrix \Rightarrow ('b::zero)\ matrix \Rightarrow ('c::zero)\ matrix$   
*combine-matrix* *f* == % *A B*. *Abs-matrix* (*combine-infmatrix* *f* (*Rep-matrix* *A*) (*Rep-matrix* *B*))

**lemma** *expand-apply-infmatrix[simp]*:  $apply\ infmatrix\ f\ A\ j\ i = f\ (A\ j\ i)$   
 <proof>

**lemma** *expand-combine-infmatrix[simp]*:  $combine\ infmatrix\ f\ A\ B\ j\ i = f\ (A\ j\ i)\ (B\ j\ i)$   
 <proof>

**constdefs**

```

commutative :: ('a => 'a => 'b) => bool
commutative f == ! x y. f x y = f y x
associative :: ('a => 'a => 'a) => bool
associative f == ! x y z. f (f x y) z = f x (f y z)

```

To reason about associativity and commutativity of operations on matrices, let's take a step back and look at the general situation: Assume that we have sets  $A$  and  $B$  with  $B \subset A$  and an abstraction  $u : A \rightarrow B$ . This abstraction has to fulfill  $u(b) = b$  for all  $b \in B$ , but is arbitrary otherwise. Each function  $f : A \times A \rightarrow A$  now induces a function  $f' : B \times B \rightarrow B$  by  $f' = u \circ f$ . It is obvious that commutativity of  $f$  implies commutativity of  $f'$ :  $f'xy = u(fxy) = u(fyx) = f'yx$ .

**lemma** *combine-infmatrix-commute*:

```

commutative f ==> commutative (combine-infmatrix f)
<proof>

```

**lemma** *combine-matrix-commute*:

```

commutative f ==> commutative (combine-matrix f)
<proof>

```

On the contrary, given an associative function  $f$  we cannot expect  $f'$  to be associative. A counterexample is given by  $A = \mathbb{Z}$ ,  $B = \{-1, 0, 1\}$ , as  $f$  we take addition on  $\mathbb{Z}$ , which is clearly associative. The abstraction is given by  $u(a) = 0$  for  $a \notin B$ . Then we have

$$f'(f'11) - 1 = u(f(u(f11)) - 1) = u(f(u2) - 1) = u(f0 - 1) = -1,$$

but on the other hand we have

$$f'1(f'1 - 1) = u(f1(u(f1 - 1))) = u(f10) = 1.$$

A way out of this problem is to assume that  $f(A \times A) \subset A$  holds, and this is what we are going to do:

**lemma** *nonzero-positions-combine-infmatrix[simp]*:  $f\ 0\ 0 = 0 \implies \text{nonzero-positions } (\text{combine-infmatrix } f\ A\ B) \subseteq (\text{nonzero-positions } A) \cup (\text{nonzero-positions } B)$   
<proof>

**lemma** *finite-nonzero-positions-Rep[simp]*:  $\text{finite } (\text{nonzero-positions } (\text{Rep-matrix } A))$   
<proof>

**lemma** *combine-infmatrix-closed [simp]*:

```

f 0 0 = 0 ==> Rep-matrix (Abs-matrix (combine-infmatrix f (Rep-matrix A)
(Rep-matrix B))) = combine-infmatrix f (Rep-matrix A) (Rep-matrix B)
<proof>

```

We need the next two lemmas only later, but it is analog to the above one, so we prove them now:

**lemma** *nonzero-positions-apply-infmatrix*[simp]:  $f \ 0 = 0 \implies \text{nonzero-positions} (\text{apply-infmatrix } f \ A) \subseteq \text{nonzero-positions } A$   
 ⟨proof⟩

**lemma** *apply-infmatrix-closed* [simp]:  
 $f \ 0 = 0 \implies \text{Rep-matrix } (\text{Abs-matrix } (\text{apply-infmatrix } f \ (\text{Rep-matrix } A))) = \text{apply-infmatrix } f \ (\text{Rep-matrix } A)$   
 ⟨proof⟩

**lemma** *combine-infmatrix-assoc*[simp]:  $f \ 0 \ 0 = 0 \implies \text{associative } f \implies \text{associative} (\text{combine-infmatrix } f)$   
 ⟨proof⟩

**lemma** *comb*:  $f = g \implies x = y \implies f \ x = g \ y$   
 ⟨proof⟩

**lemma** *combine-matrix-assoc*:  $f \ 0 \ 0 = 0 \implies \text{associative } f \implies \text{associative} (\text{combine-matrix } f)$   
 ⟨proof⟩

**lemma** *Rep-apply-matrix*[simp]:  $f \ 0 = 0 \implies \text{Rep-matrix } (\text{apply-matrix } f \ A) \ j \ i = f \ (\text{Rep-matrix } A \ j \ i)$   
 ⟨proof⟩

**lemma** *Rep-combine-matrix*[simp]:  $f \ 0 \ 0 = 0 \implies \text{Rep-matrix } (\text{combine-matrix } f \ A \ B) \ j \ i = f \ (\text{Rep-matrix } A \ j \ i) \ (\text{Rep-matrix } B \ j \ i)$   
 ⟨proof⟩

**lemma** *combine-nrows*:  $f \ 0 \ 0 = 0 \implies \text{nrows} (\text{combine-matrix } f \ A \ B) \leq \max (\text{nrows } A) (\text{nrows } B)$   
 ⟨proof⟩

**lemma** *combine-ncols*:  $f \ 0 \ 0 = 0 \implies \text{ncols} (\text{combine-matrix } f \ A \ B) \leq \max (\text{ncols } A) (\text{ncols } B)$   
 ⟨proof⟩

**lemma** *combine-nrows*:  $f \ 0 \ 0 = 0 \implies \text{nrows } A \leq q \implies \text{nrows } B \leq q \implies \text{nrows} (\text{combine-matrix } f \ A \ B) \leq q$   
 ⟨proof⟩

**lemma** *combine-ncols*:  $f \ 0 \ 0 = 0 \implies \text{ncols } A \leq q \implies \text{ncols } B \leq q \implies \text{ncols} (\text{combine-matrix } f \ A \ B) \leq q$   
 ⟨proof⟩

**constdefs**

*zero-r-neutral* ::  $('a \Rightarrow 'b :: \text{zero} \Rightarrow 'a) \Rightarrow \text{bool}$   
*zero-r-neutral*  $f == ! a. f \ a \ 0 = a$

*zero-l-neutral* :: ('a::zero ⇒ 'b ⇒ 'b) ⇒ bool  
*zero-l-neutral* f == ! a. f 0 a = a  
*zero-closed* :: (('a::zero) ⇒ ('b::zero) ⇒ ('c::zero)) ⇒ bool  
*zero-closed* f == (!x. f x 0 = 0) & (!y. f 0 y = 0)

**consts** *foldseq* :: ('a ⇒ 'a ⇒ 'a) ⇒ (nat ⇒ 'a) ⇒ nat ⇒ 'a  
**primrec**

*foldseq* f s 0 = s 0  
*foldseq* f s (Suc n) = f (s 0) (foldseq f (% k. s(Suc k)) n)

**consts** *foldseq-transposed* :: ('a ⇒ 'a ⇒ 'a) ⇒ (nat ⇒ 'a) ⇒ nat ⇒ 'a  
**primrec**

*foldseq-transposed* f s 0 = s 0  
*foldseq-transposed* f s (Suc n) = f (foldseq-transposed f s n) (s (Suc n))

**lemma** *foldseq-assoc* : associative f ⇒ foldseq f = foldseq-transposed f  
 <proof>

**lemma** *foldseq-distr*: [[associative f; commutative f]] ⇒ foldseq f (% k. f (u k) (v k)) n = f (foldseq f u n) (foldseq f v n)  
 <proof>

**theorem** [[associative f; associative g; ∀ a b c d. g (f a b) (f c d) = f (g a c) (g b d); ? x y. (f x) ≠ (f y); ? x y. (g x) ≠ (g y); f x x = x; g x x = x]] ⇒ f=g | (! y. f y x = y) | (! y. g y x = y)  
 <proof>

**lemma** *foldseq-zero*:  
**assumes** fz: f 0 0 = 0 **and** sz: ! i. i ≤ n → s i = 0  
**shows** foldseq f s n = 0  
 <proof>

**lemma** *foldseq-significant-positions*:  
**assumes** p: ! i. i ≤ N → S i = T i  
**shows** foldseq f S N = foldseq f T N (**is** ?concl)  
 <proof>

**lemma** *foldseq-tail*: M ≤ N ⇒ foldseq f S N = foldseq f (% k. (if k < M then (S k) else (foldseq f (% k. S(k+M)) (N-M)))) M (**is** ?p ⇒ ?concl)  
 <proof>

**lemma** *foldseq-zerotail*:  
**assumes**  
 fz: f 0 0 = 0  
**and** sz: ! i. n ≤ i → s i = 0  
**and** nm: n ≤ m  
**shows**  
 foldseq f s n = foldseq f s m

*<proof>*

**lemma** *foldseq-zerotail2*:

**assumes** !  $x. f\ x\ 0 = x$

**and** !  $i. n < i \longrightarrow s\ i = 0$

**and**  $nm: n \leq m$

**shows**

$foldseq\ f\ s\ n = foldseq\ f\ s\ m$  (**is** ?concl)

*<proof>*

**lemma** *foldseq-zerostart*:

!  $x. f\ 0\ (f\ 0\ x) = f\ 0\ x \implies ! i. i \leq n \longrightarrow s\ i = 0 \implies foldseq\ f\ s\ (Suc\ n) = f\ 0\ (s\ (Suc\ n))$

*<proof>*

**lemma** *foldseq-zerostart2*:

!  $x. f\ 0\ x = x \implies ! i. i < n \longrightarrow s\ i = 0 \implies foldseq\ f\ s\ n = s\ n$

*<proof>*

**lemma** *foldseq-almostzero*:

**assumes**  $f0x: ! x. f\ 0\ x = x$  **and**  $fx0: ! x. f\ x\ 0 = x$  **and**  $s0: ! i. i \neq j \longrightarrow s\ i = 0$

**shows**  $foldseq\ f\ s\ n = (if\ (j \leq n)\ then\ (s\ j)\ else\ 0)$  (**is** ?concl)

*<proof>*

**lemma** *foldseq-distr-unary*:

**assumes** !!  $a\ b. g\ (f\ a\ b) = f\ (g\ a)\ (g\ b)$

**shows**  $g\ (foldseq\ f\ s\ n) = foldseq\ f\ (\% x. g\ (s\ x))\ n$  (**is** ?concl)

*<proof>*

**constdefs**

$mult\ matrix\ n :: nat \Rightarrow (('a::zero) \Rightarrow ('b::zero) \Rightarrow ('c::zero)) \Rightarrow ('c \Rightarrow 'c \Rightarrow 'c) \Rightarrow 'a\ matrix \Rightarrow 'b\ matrix \Rightarrow 'c\ matrix$

$mult\ matrix\ n\ n\ fmul\ fadd\ A\ B == Abs\ matrix\ (\% j\ i. foldseq\ fadd\ (\% k. fmul\ (Rep\ matrix\ A\ j\ k)\ (Rep\ matrix\ B\ k\ i))\ n)$

$mult\ matrix :: (('a::zero) \Rightarrow ('b::zero) \Rightarrow ('c::zero)) \Rightarrow ('c \Rightarrow 'c \Rightarrow 'c) \Rightarrow 'a\ matrix \Rightarrow 'b\ matrix \Rightarrow 'c\ matrix$

$mult\ matrix\ fmul\ fadd\ A\ B == mult\ matrix\ n\ (max\ (ncols\ A)\ (nrows\ B))\ fmul\ fadd\ A\ B$

**lemma** *mult-matrix-n*:

**assumes**  $prems: ncols\ A \leq n$  (**is** ?An)  $nrows\ B \leq n$  (**is** ?Bn)  $fadd\ 0\ 0 = 0$   $fmul\ 0\ 0 = 0$

**shows**  $c: mult\ matrix\ fmul\ fadd\ A\ B = mult\ matrix\ n\ n\ fmul\ fadd\ A\ B$  (**is** ?concl)

*<proof>*

**lemma** *mult-matrix-nm*:

**assumes**  $prems: ncols\ A \leq n$   $nrows\ B \leq n$   $ncols\ A \leq m$   $nrows\ B \leq m$   $fadd\ 0\ 0 = 0$   $fmul\ 0\ 0 = 0$

**shows**  $mult\ matrix\ n\ n\ fmul\ fadd\ A\ B = mult\ matrix\ n\ m\ fmul\ fadd\ A\ B$

*<proof>*

**constdefs**

*r-distributive* :: ('a ⇒ 'b ⇒ 'b) ⇒ ('b ⇒ 'b ⇒ 'b) ⇒ bool  
*r-distributive fmul fadd* == ! a u v. fmul a (fadd u v) = fadd (fmul a u) (fmul a v)  
*l-distributive* :: ('a ⇒ 'b ⇒ 'a) ⇒ ('a ⇒ 'a ⇒ 'a) ⇒ bool  
*l-distributive fmul fadd* == ! a u v. fmul (fadd u v) a = fadd (fmul u a) (fmul v a)  
*distributive* :: ('a ⇒ 'a ⇒ 'a) ⇒ ('a ⇒ 'a ⇒ 'a) ⇒ bool  
*distributive fmul fadd* == *l-distributive fmul fadd* & *r-distributive fmul fadd*

**lemma** *max1*: !! a x y. (a::nat) <= x ⇒ a <= max x y *<proof>*

**lemma** *max2*: !! b x y. (b::nat) <= y ⇒ b <= max x y *<proof>*

**lemma** *r-distributive-matrix*:

**assumes** *prems*:

*r-distributive fmul fadd*

*associative fadd*

*commutative fadd*

*fadd 0 0 = 0*

! a. *fmul a 0 = 0*

! a. *fmul 0 a = 0*

**shows** *r-distributive (mult-matrix fmul fadd) (combine-matrix fadd) (is ?concl)*  
*<proof>*

**lemma** *l-distributive-matrix*:

**assumes** *prems*:

*l-distributive fmul fadd*

*associative fadd*

*commutative fadd*

*fadd 0 0 = 0*

! a. *fmul a 0 = 0*

! a. *fmul 0 a = 0*

**shows** *l-distributive (mult-matrix fmul fadd) (combine-matrix fadd) (is ?concl)*  
*<proof>*

**instance** *matrix* :: (zero) zero *<proof>*

**defs(overloaded)**

*zero-matrix-def*: (0::('a::zero) matrix) == *Abs-matrix*(% j i. 0)

**lemma** *Rep-zero-matrix-def[simp]*: *Rep-matrix 0 j i = 0*  
*<proof>*

**lemma** *zero-matrix-def-nrows[simp]*: *nrows 0 = 0*  
*<proof>*

**lemma** *zero-matrix-def-ncols[simp]*: *ncols 0 = 0*

*<proof>*

**lemma** *combine-matrix-zero-l-neutral*:  $\text{zero-l-neutral } f \implies \text{zero-l-neutral } (\text{combine-matrix } f)$

*<proof>*

**lemma** *combine-matrix-zero-r-neutral*:  $\text{zero-r-neutral } f \implies \text{zero-r-neutral } (\text{combine-matrix } f)$

*<proof>*

**lemma** *mult-matrix-zero-closed*:  $\llbracket \text{fadd } 0 \ 0 = 0; \text{zero-closed } \text{fmul} \rrbracket \implies \text{zero-closed } (\text{mult-matrix } \text{fmul } \text{fadd})$

*<proof>*

**lemma** *mult-matrix-n-zero-right[simp]*:  $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } a \ 0 = 0 \rrbracket \implies \text{mult-matrix-n } n \ \text{fmul } \text{fadd } A \ 0 = 0$

*<proof>*

**lemma** *mult-matrix-n-zero-left[simp]*:  $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } 0 \ a = 0 \rrbracket \implies \text{mult-matrix-n } n \ \text{fmul } \text{fadd } 0 \ A = 0$

*<proof>*

**lemma** *mult-matrix-zero-left[simp]*:  $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } 0 \ a = 0 \rrbracket \implies \text{mult-matrix } \text{fmul } \text{fadd } 0 \ A = 0$

*<proof>*

**lemma** *mult-matrix-zero-right[simp]*:  $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } a \ 0 = 0 \rrbracket \implies \text{mult-matrix } \text{fmul } \text{fadd } A \ 0 = 0$

*<proof>*

**lemma** *apply-matrix-zero[simp]*:  $f \ 0 = 0 \implies \text{apply-matrix } f \ 0 = 0$

*<proof>*

**lemma** *combine-matrix-zero*:  $f \ 0 \ 0 = 0 \implies \text{combine-matrix } f \ 0 \ 0 = 0$

*<proof>*

**lemma** *transpose-matrix-zero[simp]*:  $\text{transpose-matrix } 0 = 0$

*<proof>*

**lemma** *apply-zero-matrix-def[simp]*:  $\text{apply-matrix } (\% \ x. \ 0) \ A = 0$

*<proof>*

**constdefs**

*singleton-matrix* ::  $\text{nat} \Rightarrow \text{nat} \Rightarrow ('a::\text{zero}) \Rightarrow 'a \ \text{matrix}$   
*singleton-matrix*  $j \ i \ a == \text{Abs-matrix } (\% \ m \ n. \ \text{if } j = m \ \& \ i = n \ \text{then } a \ \text{else } 0)$   
*move-matrix* ::  $('a::\text{zero}) \ \text{matrix} \Rightarrow \text{int} \Rightarrow \text{int} \Rightarrow 'a \ \text{matrix}$   
*move-matrix*  $A \ y \ x == \text{Abs-matrix } (\% \ j \ i. \ \text{if } (\text{neg } ((\text{int } j) - y)) \mid (\text{neg } ((\text{int } i) - x)) \ \text{then } 0 \ \text{else } \text{Rep-matrix } A \ (\text{nat } ((\text{int } j) - y)) \ (\text{nat } ((\text{int } i) - x)))$   
*take-rows* ::  $('a::\text{zero}) \ \text{matrix} \Rightarrow \text{nat} \Rightarrow 'a \ \text{matrix}$

$take\_rows\ A\ r == Abs\_matrix(\% j\ i.\ if\ (j < r)\ then\ (Rep\_matrix\ A\ j\ i)\ else\ 0)$   
 $take\_columns :: ('a::zero)\ matrix \Rightarrow nat \Rightarrow 'a\ matrix$   
 $take\_columns\ A\ c == Abs\_matrix(\% j\ i.\ if\ (i < c)\ then\ (Rep\_matrix\ A\ j\ i)\ else\ 0)$

**constdefs**

$column\_of\_matrix :: ('a::zero)\ matrix \Rightarrow nat \Rightarrow 'a\ matrix$   
 $column\_of\_matrix\ A\ n == take\_columns\ (move\_matrix\ A\ 0\ (-\ int\ n))\ 1$   
 $row\_of\_matrix :: ('a::zero)\ matrix \Rightarrow nat \Rightarrow 'a\ matrix$   
 $row\_of\_matrix\ A\ m == take\_rows\ (move\_matrix\ A\ (-\ int\ m)\ 0)\ 1$

**lemma**  $Rep\_singleton\_matrix[simp]: Rep\_matrix\ (singleton\_matrix\ j\ i\ e)\ m\ n = (if\ j = m\ \&\ i = n\ then\ e\ else\ 0)$   
 $\langle proof \rangle$

**lemma**  $apply\_singleton\_matrix[simp]: f\ 0 = 0 \implies apply\_matrix\ f\ (singleton\_matrix\ j\ i\ x) = (singleton\_matrix\ j\ i\ (f\ x))$   
 $\langle proof \rangle$

**lemma**  $singleton\_matrix\_zero[simp]: singleton\_matrix\ j\ i\ 0 = 0$   
 $\langle proof \rangle$

**lemma**  $nrows\_singleton[simp]: nrows(singleton\_matrix\ j\ i\ e) = (if\ e = 0\ then\ 0\ else\ Suc\ j)$   
 $\langle proof \rangle$

**lemma**  $ncols\_singleton[simp]: ncols(singleton\_matrix\ j\ i\ e) = (if\ e = 0\ then\ 0\ else\ Suc\ i)$   
 $\langle proof \rangle$

**lemma**  $combine\_singleton: f\ 0\ 0 = 0 \implies combine\_matrix\ f\ (singleton\_matrix\ j\ i\ a)\ (singleton\_matrix\ j\ i\ b) = singleton\_matrix\ j\ i\ (f\ a\ b)$   
 $\langle proof \rangle$

**lemma**  $transpose\_singleton[simp]: transpose\_matrix\ (singleton\_matrix\ j\ i\ a) = singleton\_matrix\ i\ j\ a$   
 $\langle proof \rangle$

**lemma**  $Rep\_move\_matrix[simp]:$   
 $Rep\_matrix\ (move\_matrix\ A\ y\ x)\ j\ i =$   
 $(if\ (neg\ ((int\ j) - y))\ |\ (neg\ ((int\ i) - x))\ then\ 0\ else\ Rep\_matrix\ A\ (nat\ ((int\ j) - y))\ (nat\ ((int\ i) - x)))$   
 $\langle proof \rangle$

**lemma**  $move\_matrix\_0\_0[simp]: move\_matrix\ A\ 0\ 0 = A$   
 $\langle proof \rangle$

**lemma**  $move\_matrix\_ortho: move\_matrix\ A\ j\ i = move\_matrix\ (move\_matrix\ A\ j\ 0)\ 0\ i$

*<proof>*

**lemma** *transpose-move-matrix[simp]:*

*transpose-matrix (move-matrix A x y) = move-matrix (transpose-matrix A) y x*  
*<proof>*

**lemma** *move-matrix-singleton[simp]:* *move-matrix (singleton-matrix u v x) j i =*  
*(if (j + int u < 0) | (i + int v < 0) then 0 else (singleton-matrix (nat (j + int*  
*u)) (nat (i + int v)) x))*  
*<proof>*

**lemma** *Rep-take-columns[simp]:*

*Rep-matrix (take-columns A c) j i =*  
*(if i < c then (Rep-matrix A j i) else 0)*  
*<proof>*

**lemma** *Rep-take-rows[simp]:*

*Rep-matrix (take-rows A r) j i =*  
*(if j < r then (Rep-matrix A j i) else 0)*  
*<proof>*

**lemma** *Rep-column-of-matrix[simp]:*

*Rep-matrix (column-of-matrix A c) j i = (if i = 0 then (Rep-matrix A j c) else*  
*0)*  
*<proof>*

**lemma** *Rep-row-of-matrix[simp]:*

*Rep-matrix (row-of-matrix A r) j i = (if j = 0 then (Rep-matrix A r i) else 0)*  
*<proof>*

**lemma** *column-of-matrix: ncols A <= n ==> column-of-matrix A n = 0*

*<proof>*

**lemma** *row-of-matrix: nrows A <= n ==> row-of-matrix A n = 0*

*<proof>*

**lemma** *mult-matrix-singleton-right[simp]:*

**assumes** *prems:*

*! x. fmul x 0 = 0*

*! x. fmul 0 x = 0*

*! x. fadd 0 x = x*

*! x. fadd x 0 = x*

**shows** *(mult-matrix fmul fadd A (singleton-matrix j i e)) = apply-matrix (% x.*  
*fmul x e) (move-matrix (column-of-matrix A j) 0 (int i))*

*<proof>*

**lemma** *mult-matrix-ext:*

**assumes**

*eprem:*

```

? e. (! a b. a ≠ b → fmul a e ≠ fmul b e)
and fprems:
! a. fmul 0 a = 0
! a. fmul a 0 = 0
! a. fadd a 0 = a
! a. fadd 0 a = a
and contraprems:
mult-matrix fmul fadd A = mult-matrix fmul fadd B
shows
A = B
⟨proof⟩

constdefs
foldmatrix :: ('a ⇒ 'a ⇒ 'a) ⇒ ('a ⇒ 'a ⇒ 'a) ⇒ ('a infmatrix) ⇒ nat ⇒ nat
⇒ 'a
foldmatrix f g A m n == foldseq-transposed g (% j. foldseq f (A j) n) m
foldmatrix-transposed :: ('a ⇒ 'a ⇒ 'a) ⇒ ('a ⇒ 'a ⇒ 'a) ⇒ ('a infmatrix) ⇒
nat ⇒ nat ⇒ 'a
foldmatrix-transposed f g A m n == foldseq g (% j. foldseq-transposed f (A j) n)
m

lemma foldmatrix-transpose:
assumes
! a b c d. g(f a b) (f c d) = f (g a c) (g b d)
shows
foldmatrix f g A m n = foldmatrix-transposed g f (transpose-infmatrix A) n m
(is ?concl)
⟨proof⟩

lemma foldseq-foldseq:
assumes
associative f
associative g
! a b c d. g(f a b) (f c d) = f (g a c) (g b d)
shows
foldseq g (% j. foldseq f (A j) n) m = foldseq f (% j. foldseq g ((transpose-infmatrix
A) j) m) n
⟨proof⟩

lemma mult-n-nrows:
assumes
! a. fmul 0 a = 0
! a. fmul a 0 = 0
fadd 0 0 = 0
shows nrows (mult-matrix-n n fmul fadd A B) ≤ nrows A
⟨proof⟩

lemma mult-n-ncols:
assumes

```

! a.  $fmul\ 0\ a = 0$   
 ! a.  $fmul\ a\ 0 = 0$   
 $fadd\ 0\ 0 = 0$   
**shows**  $ncols\ (mult\ matrix\ n\ n\ fmul\ fadd\ A\ B) \leq ncols\ B$   
 <proof>

**lemma** *mult-nrows:*

**assumes**

! a.  $fmul\ 0\ a = 0$

! a.  $fmul\ a\ 0 = 0$

$fadd\ 0\ 0 = 0$

**shows**  $nrows\ (mult\ matrix\ fmul\ fadd\ A\ B) \leq nrows\ A$

<proof>

**lemma** *mult-ncols:*

**assumes**

! a.  $fmul\ 0\ a = 0$

! a.  $fmul\ a\ 0 = 0$

$fadd\ 0\ 0 = 0$

**shows**  $ncols\ (mult\ matrix\ fmul\ fadd\ A\ B) \leq ncols\ B$

<proof>

**lemma** *nrows-move-matrix-le:*  $nrows\ (move\ matrix\ A\ j\ i) \leq nat((int\ (nrows\ A)) + j)$

<proof>

**lemma** *ncols-move-matrix-le:*  $ncols\ (move\ matrix\ A\ j\ i) \leq nat((int\ (ncols\ A)) + i)$

<proof>

**lemma** *mult-matrix-assoc:*

**assumes** *prems:*

! a.  $fmul1\ 0\ a = 0$

! a.  $fmul1\ a\ 0 = 0$

! a.  $fmul2\ 0\ a = 0$

! a.  $fmul2\ a\ 0 = 0$

$fadd1\ 0\ 0 = 0$

$fadd2\ 0\ 0 = 0$

! a b c d.  $fadd2\ (fadd1\ a\ b)\ (fadd1\ c\ d) = fadd1\ (fadd2\ a\ c)\ (fadd2\ b\ d)$

*associative fadd1*

*associative fadd2*

! a b c.  $fmul2\ (fmul1\ a\ b)\ c = fmul1\ a\ (fmul2\ b\ c)$

! a b c.  $fmul2\ (fadd1\ a\ b)\ c = fadd1\ (fmul2\ a\ c)\ (fmul2\ b\ c)$

! a b c.  $fmul1\ c\ (fadd2\ a\ b) = fadd2\ (fmul1\ c\ a)\ (fmul1\ c\ b)$

**shows**  $mult\ matrix\ fmul2\ fadd2\ (mult\ matrix\ fmul1\ fadd1\ A\ B)\ C = mult\ matrix\ fmul1\ fadd1\ A\ (mult\ matrix\ fmul2\ fadd2\ B\ C)$  (**is** ?concl)

<proof>

**lemma**

**assumes** *prems*:  
! a. *fmul1* 0 a = 0  
! a. *fmul1* a 0 = 0  
! a. *fmul2* 0 a = 0  
! a. *fmul2* a 0 = 0  
*fadd1* 0 0 = 0  
*fadd2* 0 0 = 0  
! a b c d. *fadd2* (*fadd1* a b) (*fadd1* c d) = *fadd1* (*fadd2* a c) (*fadd2* b d)  
*associative fadd1*  
*associative fadd2*  
! a b c. *fmul2* (*fmul1* a b) c = *fmul1* a (*fmul2* b c)  
! a b c. *fmul2* (*fadd1* a b) c = *fadd1* (*fmul2* a c) (*fmul2* b c)  
! a b c. *fmul1* c (*fadd2* a b) = *fadd2* (*fmul1* c a) (*fmul1* c b)  
**shows**  
(*mult-matrix fmul1 fadd1* A) o (*mult-matrix fmul2 fadd2* B) = *mult-matrix fmul2 fadd2* (*mult-matrix fmul1 fadd1* A B)  
⟨*proof*⟩

**lemma** *mult-matrix-assoc-simple*:

**assumes** *prems*:  
! a. *fmul* 0 a = 0  
! a. *fmul* a 0 = 0  
*fadd* 0 0 = 0  
*associative fadd*  
*commutative fadd*  
*associative fmul*  
*distributive fmul fadd*  
**shows** *mult-matrix fmul fadd* (*mult-matrix fmul fadd* A B) C = *mult-matrix fmul fadd* A (*mult-matrix fmul fadd* B C) (**is** ?*concl*)  
⟨*proof*⟩

**lemma** *transpose-apply-matrix*:  $f\ 0 = 0 \implies \text{transpose-matrix } (\text{apply-matrix } f\ A) = \text{apply-matrix } f\ (\text{transpose-matrix } A)$   
⟨*proof*⟩

**lemma** *transpose-combine-matrix*:  $f\ 0\ 0 = 0 \implies \text{transpose-matrix } (\text{combine-matrix } f\ A\ B) = \text{combine-matrix } f\ (\text{transpose-matrix } A)\ (\text{transpose-matrix } B)$   
⟨*proof*⟩

**lemma** *Rep-mult-matrix*:

**assumes**  
! a. *fmul* 0 a = 0  
! a. *fmul* a 0 = 0  
*fadd* 0 0 = 0  
**shows**  
*Rep-matrix*(*mult-matrix fmul fadd* A B) j i =  
*foldseq fadd* (% k. *fmul* (*Rep-matrix* A j k) (*Rep-matrix* B k i)) (max (ncols A) (nrows B))  
⟨*proof*⟩

**lemma** *transpose-mult-matrix*:

**assumes**

! a. *fmul* 0 a = 0

! a. *fmul* a 0 = 0

*fadd* 0 0 = 0

! x y. *fmul* y x = *fmul* x y

**shows**

*transpose-matrix* (*mult-matrix* *fmul* *fadd* A B) = *mult-matrix* *fmul* *fadd* (*transpose-matrix* B) (*transpose-matrix* A)

*<proof>*

**lemma** *column-transpose-matrix*: *column-of-matrix* (*transpose-matrix* A) n = *transpose-matrix* (*row-of-matrix* A n)

*<proof>*

**lemma** *take-columns-transpose-matrix*: *take-columns* (*transpose-matrix* A) n = *transpose-matrix* (*take-rows* A n)

*<proof>*

**instance** *matrix* :: ({ord, zero}) ord *<proof>*

**defs** (overloaded)

*le-matrix-def*: (A::('a::{ord,zero}) *matrix*) <= B == ! j i. (*Rep-matrix* A j i) <= (*Rep-matrix* B j i)

*less-def*: (A::('a::{ord,zero}) *matrix*) < B == (A <= B) & (A ≠ B)

**instance** *matrix* :: ({order, zero}) order

*<proof>*

**lemma** *le-apply-matrix*:

**assumes**

f 0 = 0

! x y. x <= y → f x <= f y

(a::('a::{ord, zero}) *matrix*) <= b

**shows**

*apply-matrix* f a <= *apply-matrix* f b

*<proof>*

**lemma** *le-combine-matrix*:

**assumes**

f 0 0 = 0

! a b c d. a <= b & c <= d → f a c <= f b d

A <= B

C <= D

**shows**

*combine-matrix* f A C <= *combine-matrix* f B D

*<proof>*

**lemma** *le-left-combine-matrix*:

**assumes**

$f\ 0\ 0 = 0$

$! a\ b\ c. a \leq b \longrightarrow f\ c\ a \leq f\ c\ b$

$A \leq B$

**shows**

$combine\ matrix\ f\ C\ A \leq combine\ matrix\ f\ C\ B$

*<proof>*

**lemma** *le-right-combine-matrix*:

**assumes**

$f\ 0\ 0 = 0$

$! a\ b\ c. a \leq b \longrightarrow f\ a\ c \leq f\ b\ c$

$A \leq B$

**shows**

$combine\ matrix\ f\ A\ C \leq combine\ matrix\ f\ B\ C$

*<proof>*

**lemma** *le-transpose-matrix*:  $(A \leq B) = (transpose\ matrix\ A \leq transpose\ matrix\ B)$

*<proof>*

**lemma** *le-foldseq*:

**assumes**

$! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow f\ a\ c \leq f\ b\ d$

$! i. i \leq n \longrightarrow s\ i \leq t\ i$

**shows**

$foldseq\ f\ s\ n \leq foldseq\ f\ t\ n$

*<proof>*

**lemma** *le-left-mult*:

**assumes**

$! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow fadd\ a\ c \leq fadd\ b\ d$

$! c\ a\ b. 0 \leq c \ \& \ a \leq b \longrightarrow fmul\ c\ a \leq fmul\ c\ b$

$! a. fmul\ 0\ a = 0$

$! a. fmul\ a\ 0 = 0$

$fadd\ 0\ 0 = 0$

$0 \leq C$

$A \leq B$

**shows**

$mult\ matrix\ fmul\ fadd\ C\ A \leq mult\ matrix\ fmul\ fadd\ C\ B$

*<proof>*

**lemma** *le-right-mult*:

**assumes**

$! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow fadd\ a\ c \leq fadd\ b\ d$

$! c\ a\ b. 0 \leq c \ \& \ a \leq b \longrightarrow fmul\ a\ c \leq fmul\ b\ c$

$! a. fmul\ 0\ a = 0$

$! a. fmul\ a\ 0 = 0$

```

fadd 0 0 = 0
0 <= C
A <= B
shows
mult-matrix fmul fadd A C <= mult-matrix fmul fadd B C
⟨proof⟩

lemma spec2: ! j i. P j i ⟹ P j i ⟨proof⟩
lemma neg-imp: (¬ Q ⟶ ¬ P) ⟹ P ⟶ Q ⟨proof⟩

lemma singleton-matrix-le[simp]: (singleton-matrix j i a <= singleton-matrix j i
b) = (a <= (b:::order))
⟨proof⟩

lemma singleton-le-zero[simp]: (singleton-matrix j i x <= 0) = (x <= (0::'a::{order,zero}))
⟨proof⟩

lemma singleton-ge-zero[simp]: (0 <= singleton-matrix j i x) = ((0::'a::{order,zero})
<= x)
⟨proof⟩

lemma move-matrix-le-zero[simp]: 0 <= j ⟹ 0 <= i ⟹ (move-matrix A j i
<= 0) = (A <= (0::('a::{order,zero}) matrix))
⟨proof⟩

lemma move-matrix-zero-le[simp]: 0 <= j ⟹ 0 <= i ⟹ (0 <= move-matrix
A j i) = ((0::('a::{order,zero}) matrix) <= A)
⟨proof⟩

lemma move-matrix-le-move-matrix-iff[simp]: 0 <= j ⟹ 0 <= i ⟹ (move-matrix
A j i <= move-matrix B j i) = (A <= (B::('a::{order,zero}) matrix))
⟨proof⟩

end

theory Matrix=MatrixGeneral:

instance matrix :: (minus) minus
⟨proof⟩

instance matrix :: (plus) plus
⟨proof⟩

instance matrix :: ({plus,times}) times
⟨proof⟩

defs (overloaded)

```

*plus-matrix-def*:  $A + B == \text{combine-matrix } (op +) A B$   
*diff-matrix-def*:  $A - B == \text{combine-matrix } (op -) A B$   
*minus-matrix-def*:  $- A == \text{apply-matrix } \text{uminus } A$   
*times-matrix-def*:  $A * B == \text{mult-matrix } (op *) (op +) A B$

**lemma** *is-meet-combine-matrix-meet*: *is-meet* (*combine-matrix meet*)  
 ⟨*proof*⟩

**lemma** *is-join-combine-matrix-join*: *is-join* (*combine-matrix join*)  
 ⟨*proof*⟩

**instance** *matrix* :: (*lordered-ab-group*) *lordered-ab-group-meet*  
 ⟨*proof*⟩

**defs** (**overloaded**)  
*abs-matrix-def*: *abs* ( $A :: ('a :: \text{lordered-ab-group}) \text{matrix}$ ) == *join*  $A (- A)$

**instance** *matrix* :: (*lordered-ring*) *lordered-ring*  
 ⟨*proof*⟩

**lemma** *Rep-matrix-add*[*simp*]: *Rep-matrix* ( $((a :: ('a :: \text{lordered-ab-group}) \text{matrix}) + b)$   
 $j i = (\text{Rep-matrix } a j i) + (\text{Rep-matrix } b j i)$   
 ⟨*proof*⟩

**lemma** *Rep-matrix-mult*: *Rep-matrix* ( $((a :: ('a :: \text{lordered-ring}) \text{matrix}) * b)$   $j i =$   
 $\text{foldseq } (op +) (\% k. (\text{Rep-matrix } a j k) * (\text{Rep-matrix } b k i)) (\text{max } (nrows a)$   
 $(nrows b))$   
 ⟨*proof*⟩

**lemma** *apply-matrix-add*:  $! x y. f (x+y) = (f x) + (f y) \implies f 0 = (0 :: 'a) \implies$   
 $\text{apply-matrix } f ((a :: ('a :: \text{lordered-ab-group}) \text{matrix}) + b) = (\text{apply-matrix } f a) +$   
 $(\text{apply-matrix } f b)$   
 ⟨*proof*⟩

**lemma** *singleton-matrix-add*: *singleton-matrix*  $j i ((a :: (\text{lordered-ab-group}) + b) =$   
 $(\text{singleton-matrix } j i a) + (\text{singleton-matrix } j i b)$   
 ⟨*proof*⟩

**lemma** *nrows-mult*:  $nrows ((A :: ('a :: \text{lordered-ring}) \text{matrix}) * B) <= nrows A$   
 ⟨*proof*⟩

**lemma** *ncols-mult*:  $ncols ((A :: ('a :: \text{lordered-ring}) \text{matrix}) * B) <= ncols B$   
 ⟨*proof*⟩

**constdefs**  
*one-matrix* ::  $\text{nat} \Rightarrow ('a :: \{\text{zero, one}\}) \text{matrix}$   
*one-matrix*  $n == \text{Abs-matrix } (\% j i. \text{if } j = i \ \& \ j < n \ \text{then } 1 \ \text{else } 0)$

**lemma** *Rep-one-matrix*[simp]: *Rep-matrix (one-matrix n) j i = (if (j = i & j < n) then 1 else 0)*

*<proof>*

**lemma** *nrows-one-matrix*[simp]: *nrows ((one-matrix n) :: ('a::axclass-0-neq-1)matrix) = n (is ?r = -)*

*<proof>*

**lemma** *ncols-one-matrix*[simp]: *ncols ((one-matrix n) :: ('a::axclass-0-neq-1)matrix) = n (is ?r = -)*

*<proof>*

**lemma** *one-matrix-mult-right*[simp]: *ncols A <= n  $\implies$  (A::('a::{lordered-ring, ring-1}) matrix) \* (one-matrix n) = A*

*<proof>*

**lemma** *one-matrix-mult-left*[simp]: *nrows A <= n  $\implies$  (one-matrix n) \* A = (A::('a::{lordered-ring, ring-1}) matrix)*

*<proof>*

**lemma** *transpose-matrix-mult*: *transpose-matrix ((A::('a::{lordered-ring, comm-ring}) matrix)\*B) = (transpose-matrix B) \* (transpose-matrix A)*

*<proof>*

**lemma** *transpose-matrix-add*: *transpose-matrix ((A::('a::lordered-ab-group) matrix)+B) = transpose-matrix A + transpose-matrix B*

*<proof>*

**lemma** *transpose-matrix-diff*: *transpose-matrix ((A::('a::lordered-ab-group) matrix)-B) = transpose-matrix A - transpose-matrix B*

*<proof>*

**lemma** *transpose-matrix-minus*: *transpose-matrix (-(A::('a::lordered-ring) matrix)) = - transpose-matrix (A::('a::lordered-ring) matrix)*

*<proof>*

### constdefs

*right-inverse-matrix* :: ('a::{lordered-ring, ring-1}) matrix  $\Rightarrow$  'a matrix  $\Rightarrow$  bool  
*right-inverse-matrix* A X == (A \* X = one-matrix (max (nrows A) (ncols X)))  
 $\wedge$  nrows X  $\leq$  ncols A

*left-inverse-matrix* :: ('a::{lordered-ring, ring-1}) matrix  $\Rightarrow$  'a matrix  $\Rightarrow$  bool  
*left-inverse-matrix* A X == (X \* A = one-matrix (max(nrows X) (ncols A)))  $\wedge$   
ncols X  $\leq$  nrows A

*inverse-matrix* :: ('a::{lordered-ring, ring-1}) matrix  $\Rightarrow$  'a matrix  $\Rightarrow$  bool  
*inverse-matrix* A X == (right-inverse-matrix A X)  $\wedge$  (left-inverse-matrix A X)

**lemma** *right-inverse-matrix-dim*: *right-inverse-matrix A X  $\implies$  nrows A = ncols X*

*<proof>*

**lemma** *left-inverse-matrix-dim*: *left-inverse-matrix*  $A$   $Y \implies \text{ncols } A = \text{nrows } Y$   
 ⟨proof⟩

**lemma** *left-right-inverse-matrix-unique*:  
 assumes *left-inverse-matrix*  $A$   $Y$  *right-inverse-matrix*  $A$   $X$   
 shows  $X = Y$   
 ⟨proof⟩

**lemma** *inverse-matrix-inject*:  $\llbracket \text{inverse-matrix } A$   $X; \text{inverse-matrix } A$   $Y \rrbracket \implies X = Y$   
 ⟨proof⟩

**lemma** *one-matrix-inverse*: *inverse-matrix* (*one-matrix*  $n$ ) (*one-matrix*  $n$ )  
 ⟨proof⟩

**lemma** *zero-imp-mult-zero*: ( $a::'a::\text{ring}$ ) = 0 |  $b = 0 \implies a * b = 0$   
 ⟨proof⟩

**lemma** *Rep-matrix-zero-imp-mult-zero*:  
 !  $j$   $i$   $k$ . (*Rep-matrix*  $A$   $j$   $k = 0$ ) | (*Rep-matrix*  $B$   $k$   $i$ ) = 0  $\implies A * B =$   
 ( $0::('a::\text{lordered-ring})$  *matrix*)  
 ⟨proof⟩

**lemma** *add-nrows*: *nrows* ( $A::('a::\text{comm-monoid-add})$  *matrix*)  $\leq u \implies \text{nrows } B \leq u \implies \text{nrows } (A + B) \leq u$   
 ⟨proof⟩

**lemma** *move-matrix-row-mult*: *move-matrix* (( $A::('a::\text{lordered-ring})$  *matrix*) \*  $B$ )  
 $j$   $0 = (\text{move-matrix } A$   $j$   $0) * B$   
 ⟨proof⟩

**lemma** *move-matrix-col-mult*: *move-matrix* (( $A::('a::\text{lordered-ring})$  *matrix*) \*  $B$ )  
 $0$   $i = A * (\text{move-matrix } B$   $0$   $i)$   
 ⟨proof⟩

**lemma** *move-matrix-add*: ((*move-matrix* ( $A + B$ )  $j$   $i$ )::( $'a::\text{lordered-ab-group}$ )  
*matrix*) = (*move-matrix*  $A$   $j$   $i$ ) + (*move-matrix*  $B$   $j$   $i$ )  
 ⟨proof⟩

**lemma** *move-matrix-mult*: *move-matrix* (( $A::('a::\text{lordered-ring})$  *matrix*)\* $B$ )  $j$   $i =$   
 (*move-matrix*  $A$   $j$   $0$ ) \* (*move-matrix*  $B$   $0$   $i$ )  
 ⟨proof⟩

**constdefs**

*scalar-mult* :: ( $'a::\text{lordered-ring}$ )  $\Rightarrow 'a$  *matrix*  $\Rightarrow 'a$  *matrix*  
*scalar-mult*  $a$   $m == \text{apply-matrix } (\text{op } * a) m$

**lemma** *scalar-mult-zero[simp]*: *scalar-mult*  $y$   $0 = 0$

$\langle \text{proof} \rangle$

**lemma** *scalar-mult-add*:  $\text{scalar-mult } y (a+b) = (\text{scalar-mult } y a) + (\text{scalar-mult } y b)$ 
  
 $\langle \text{proof} \rangle$

**lemma** *Rep-scalar-mult[simp]*:  $\text{Rep-matrix } (\text{scalar-mult } y a) j i = y * (\text{Rep-matrix } a j i)$ 
  
 $\langle \text{proof} \rangle$

**lemma** *scalar-mult-singleton[simp]*:  $\text{scalar-mult } y (\text{singleton-matrix } j i x) = \text{singleton-matrix } j i (y * x)$ 
  
 $\langle \text{proof} \rangle$

**lemma** *Rep-minus[simp]*:  $\text{Rep-matrix } (-(A:::\text{ordered-ab-group})) x y = - (\text{Rep-matrix } A x y)$ 
  
 $\langle \text{proof} \rangle$

**lemma** *join-matrix*:  $\text{join } (A::('a::\text{ordered-ring}) \text{ matrix}) B = \text{combine-matrix } \text{join } A B$ 
  
 $\langle \text{proof} \rangle$

**lemma** *meet-matrix*:  $\text{meet } (A::('a::\text{ordered-ring}) \text{ matrix}) B = \text{combine-matrix } \text{meet } A B$ 
  
 $\langle \text{proof} \rangle$

**lemma** *Rep-abs[simp]*:  $\text{Rep-matrix } (\text{abs } (A:::\text{ordered-ring})) x y = \text{abs } (\text{Rep-matrix } A x y)$ 
  
 $\langle \text{proof} \rangle$

**end**

**theory** *SparseMatrix* **imports** *Matrix* **begin**

**types**

$'a \text{ svec} = (\text{nat} * 'a) \text{ list}$   
 $'a \text{ smat} = ('a \text{ svec}) \text{ svec}$

**consts**

$\text{sparse-row-vector} :: ('a::\text{ordered-ring}) \text{ svec} \Rightarrow 'a \text{ matrix}$   
 $\text{sparse-row-matrix} :: ('a::\text{ordered-ring}) \text{ smat} \Rightarrow 'a \text{ matrix}$

**defs**

$\text{sparse-row-vector-def} : \text{sparse-row-vector } arr == \text{foldl } (\% m x. m + (\text{singleton-matrix } 0 (\text{fst } x) (\text{snd } x))) 0 arr$   
 $\text{sparse-row-matrix-def} : \text{sparse-row-matrix } arr == \text{foldl } (\% m r. m + (\text{move-matrix } (\text{sparse-row-vector } (\text{snd } r)) (\text{int } (\text{fst } r)) 0)) 0 arr$

**lemma** *sparse-row-vector-empty*[simp]: *sparse-row-vector* [] = 0  
 ⟨proof⟩

**lemma** *sparse-row-matrix-empty*[simp]: *sparse-row-matrix* [] = 0  
 ⟨proof⟩

**lemma** *foldl-distrstart*[rule-format]: ! *a x y*. (*f* (*g x y*) *a* = *g x* (*f y a*))  $\implies$  ! *x y*.  
 (*foldl f* (*g x y*) *l* = *g x* (*foldl f y l*))  
 ⟨proof⟩

**lemma** *sparse-row-vector-cons*[simp]: *sparse-row-vector* (*a*#*arr*) = (*singleton-matrix* 0 (*fst a*) (*snd a*)) + (*sparse-row-vector* *arr*)  
 ⟨proof⟩

**lemma** *sparse-row-vector-append*[simp]: *sparse-row-vector* (*a* @ *b*) = (*sparse-row-vector* *a*) + (*sparse-row-vector* *b*)  
 ⟨proof⟩

**lemma** *nrows-spvec*[simp]: *nrows* (*sparse-row-vector* *x*) <= (*Suc* 0)  
 ⟨proof⟩

**lemma** *sparse-row-matrix-cons*: *sparse-row-matrix* (*a*#*arr*) = ((*move-matrix* (*sparse-row-vector* (*snd a*)) (*int* (*fst a*)) 0)) + *sparse-row-matrix* *arr*  
 ⟨proof⟩

**lemma** *sparse-row-matrix-append*: *sparse-row-matrix* (*arr*@*brr*) = (*sparse-row-matrix* *arr*) + (*sparse-row-matrix* *brr*)  
 ⟨proof⟩

### consts

*sorted-spvec* :: 'a *spvec*  $\Rightarrow$  *bool*  
*sorted-spmat* :: 'a *spmat*  $\Rightarrow$  *bool*

### primrec

*sorted-spmat* [] = *True*  
*sorted-spmat* (*a*#*as*) = ((*sorted-spvec* (*snd a*)) & (*sorted-spmat* *as*))

### primrec

*sorted-spvec* [] = *True*  
*sorted-spvec-step*: *sorted-spvec* (*a*#*as*) = (*case as of* []  $\Rightarrow$  *True* | *b*#*bs*  $\Rightarrow$  ((*fst a* < *fst b*) & (*sorted-spvec* *as*)))

**declare** *sorted-spvec.simps* [simp del]

**lemma** *sorted-spvec-empty*[simp]: *sorted-spvec* [] = *True*  
 ⟨proof⟩

**lemma** *sorted-spvec-cons1*: *sorted-spvec* (*a*#*as*)  $\implies$  *sorted-spvec* *as*

*<proof>*

**lemma** *sorted-spvec-cons2*:  $\text{sorted-spvec } (a\#b\#t) \implies \text{sorted-spvec } (a\#t)$   
*<proof>*

**lemma** *sorted-spvec-cons3*:  $\text{sorted-spvec}(a\#b\#t) \implies \text{fst } a < \text{fst } b$   
*<proof>*

**lemma** *sorted-sparse-row-vector-zero*[*rule-format*]:  $m \leq n \longrightarrow \text{sorted-spvec } ((n,a)\#\text{arr})$   
 $\longrightarrow \text{Rep-matrix } (\text{sparse-row-vector } \text{arr}) \text{ } j \text{ } m = 0$   
*<proof>*

**lemma** *sorted-sparse-row-matrix-zero*[*rule-format*]:  $m \leq n \longrightarrow \text{sorted-spvec } ((n,a)\#\text{arr})$   
 $\longrightarrow \text{Rep-matrix } (\text{sparse-row-matrix } \text{arr}) \text{ } m \text{ } j = 0$   
*<proof>*

**consts**

*abs-spvec* :: (*'a::lordered-ring*) *spvec*  $\Rightarrow$  *'a spvec*  
*minus-spvec* :: (*'a::lordered-ring*) *spvec*  $\Rightarrow$  *'a spvec*  
*smult-spvec* :: (*'a::lordered-ring*)  $\Rightarrow$  *'a spvec*  $\Rightarrow$  *'a spvec*  
*addmult-spvec* :: (*'a::lordered-ring*) \* *'a spvec* \* *'a spvec*  $\Rightarrow$  *'a spvec*

**primrec**

*minus-spvec* [] = []  
*minus-spvec* (*a#as*) = (*fst a*,  $-(\text{snd } a)$ )#(*minus-spvec as*)

**primrec**

*abs-spvec* [] = []  
*abs-spvec* (*a#as*) = (*fst a*, *abs (snd a)*)#(*abs-spvec as*)

**lemma** *sparse-row-vector-minus*:

*sparse-row-vector* (*minus-spvec v*) =  $-$  (*sparse-row-vector v*)  
*<proof>*

**lemma** *sparse-row-vector-abs*:

*sorted-spvec v*  $\implies$  *sparse-row-vector* (*abs-spvec v*) = *abs* (*sparse-row-vector v*)  
*<proof>*

**lemma** *sorted-spvec-minus-spvec*:

*sorted-spvec v*  $\implies$  *sorted-spvec* (*minus-spvec v*)  
*<proof>*

**lemma** *sorted-spvec-minus-spvec*:

*sorted-spvec v*  $\implies$  *sorted-spvec* (*minus-spvec v*)  
*<proof>*

**lemma** *sorted-spvec-abs-spvec*:

*sorted-spvec v*  $\implies$  *sorted-spvec* (*abs-spvec v*)  
*<proof>*

**defs**

*smult-spvec-def*:  $smult\text{-}spvec\ y\ arr == map\ (\% a.\ (fst\ a,\ y * snd\ a))\ arr$

**lemma** *smult-spvec-empty[simp]*:  $smult\text{-}spvec\ y\ [] = []$   
(*proof*)

**lemma** *smult-spvec-cons*:  $smult\text{-}spvec\ y\ (a\#\ arr) = (fst\ a,\ y * (snd\ a))\ \#\ (smult\text{-}spvec\ y\ arr)$   
(*proof*)

**recdef** *addmult-spvec measure* (% (y, a, b). length a + (length b))  
*addmult-spvec* (y, arr, []) = arr  
*addmult-spvec* (y, [], brr) = *smult-spvec* y brr  
*addmult-spvec* (y, a#arr, b#brr) = (  
  if (fst a) < (fst b) then (a#(addmult-spvec (y, arr, b#brr)))  
  else (if (fst b < fst a) then ((fst b, y \* (snd b))#(addmult-spvec (y, a#arr, brr)))  
  else ((fst a, (snd a) + y\*(snd b))#(addmult-spvec (y, arr, brr))))))

**lemma** *addmult-spvec-empty1[simp]*:  $addmult\text{-}spvec\ (y,\ [],\ a) = smult\text{-}spvec\ y\ a$   
(*proof*)

**lemma** *addmult-spvec-empty2[simp]*:  $addmult\text{-}spvec\ (y,\ a,\ []) = a$   
(*proof*)

**lemma** *sparse-row-vector-map*:  $(! x\ y.\ f\ (x+y) = (f\ x) + (f\ y)) \implies (f::'a \Rightarrow ('a::lordered-ring))\ 0 = 0 \implies$   
 $sparse\text{-}row\text{-}vector\ (map\ (\% x.\ (fst\ x,\ f\ (snd\ x)))\ a) = apply\ matrix\ f\ (sparse\text{-}row\text{-}vector\ a)$   
(*proof*)

**lemma** *sparse-row-vector-smult*:  $sparse\text{-}row\text{-}vector\ (smult\text{-}spvec\ y\ a) = scalar\text{-}mult\ y\ (sparse\text{-}row\text{-}vector\ a)$   
(*proof*)

**lemma** *sparse-row-vector-addmult-spvec*:  $sparse\text{-}row\text{-}vector\ (addmult\text{-}spvec\ (y::'a::lordered-ring,\ a,\ b)) =$   
 $(sparse\text{-}row\text{-}vector\ a) + (scalar\text{-}mult\ y\ (sparse\text{-}row\text{-}vector\ b))$   
(*proof*)

**lemma** *sorted-smult-spvec[rule-format]*:  $sorted\text{-}spvec\ a \implies sorted\text{-}spvec\ (smult\text{-}spvec\ y\ a)$   
(*proof*)

**lemma** *sorted-spvec-addmult-spvec-helper*:  $\llbracket sorted\text{-}spvec\ (addmult\text{-}spvec\ (y,\ (a,\ b)\ \#\ arr,\ brr));\ aa < a;\ sorted\text{-}spvec\ ((a,\ b)\ \#\ arr);\ sorted\text{-}spvec\ ((aa,\ ba)\ \#\ brr)\rrbracket \implies sorted\text{-}spvec\ ((aa,\ y * ba)\ \#\ addmult\text{-}spvec\ (y,\ (a,\ b)\ \#\ arr,\ brr))$

*<proof>*

**lemma** *sorted-spvec-addmult-spvec-helper2*:

$\llbracket \text{sorted-spvec } (\text{addmult-spvec } (y, \text{arr}, (aa, ba) \# \text{brr})); a < aa; \text{sorted-spvec } ((a, b) \# \text{arr}); \text{sorted-spvec } ((aa, ba) \# \text{brr}) \rrbracket$   
 $\implies \text{sorted-spvec } ((a, b) \# \text{addmult-spvec } (y, \text{arr}, (aa, ba) \# \text{brr}))$   
*<proof>*

**lemma** *sorted-spvec-addmult-spvec-helper3*[*rule-format*]:

$\text{sorted-spvec } (\text{addmult-spvec } (y, \text{arr}, \text{brr})) \longrightarrow \text{sorted-spvec } ((aa, b) \# \text{arr}) \longrightarrow$   
 $\text{sorted-spvec } ((aa, ba) \# \text{brr})$   
 $\longrightarrow \text{sorted-spvec } ((aa, b + y * ba) \# (\text{addmult-spvec } (y, \text{arr}, \text{brr})))$   
*<proof>*

**lemma** *sorted-addmult-spvec*[*rule-format*]:  $\text{sorted-spvec } a \longrightarrow \text{sorted-spvec } b \longrightarrow$   
 $\text{sorted-spvec } (\text{addmult-spvec } (y, a, b))$

*<proof>*

**consts**

$\text{mult-spvec-spmat} :: ('a::\text{lordered-ring}) \text{ spvec} * 'a \text{ spvec} * 'a \text{ smat} \Rightarrow 'a \text{ spvec}$

**recdef** *mult-spvec-spmat measure* (% (c, arr, brr). (length arr) + (length brr))

$\text{mult-spvec-spmat } (c, [], \text{brr}) = c$   
 $\text{mult-spvec-spmat } (c, \text{arr}, []) = c$   
 $\text{mult-spvec-spmat } (c, a\#\text{arr}, b\#\text{brr}) =$   
 $\text{if } ((\text{fst } a) < (\text{fst } b)) \text{ then } (\text{mult-spvec-spmat } (c, \text{arr}, b\#\text{brr}))$   
 $\text{else } (\text{if } ((\text{fst } b) < (\text{fst } a)) \text{ then } (\text{mult-spvec-spmat } (c, a\#\text{arr}, \text{brr}))$   
 $\text{else } (\text{mult-spvec-spmat } (\text{addmult-spvec } (\text{snd } a, c, \text{snd } b), \text{arr}, \text{brr})))$

**lemma** *sparse-row-mult-spvec-spmat*[*rule-format*]:  $\text{sorted-spvec } (a::('a::\text{lordered-ring})$   
 $\text{spvec}) \longrightarrow \text{sorted-spvec } B \longrightarrow$

$\text{sparse-row-vector } (\text{mult-spvec-spmat } (c, a, B)) = (\text{sparse-row-vector } c) + (\text{sparse-row-vector } a) * (\text{sparse-row-matrix } B)$   
*<proof>*

**lemma** *sorted-mult-spvec-spmat*[*rule-format*]:

$\text{sorted-spvec } (c::('a::\text{lordered-ring}) \text{ spvec}) \longrightarrow \text{sorted-spmat } B \longrightarrow \text{sorted-spvec}$   
 $(\text{mult-spvec-spmat } (c, a, B))$   
*<proof>*

**consts**

$\text{mult-spmat} :: ('a::\text{lordered-ring}) \text{ smat} \Rightarrow 'a \text{ smat} \Rightarrow 'a \text{ smat}$

**primrec**

$\text{mult-spmat } [] A = []$   
 $\text{mult-spmat } (a\#as) A = (\text{fst } a, \text{mult-spmat } ([], \text{snd } a, A))\#(\text{mult-spmat } as$   
 $A)$

**lemma** *sparse-row-mult-spmat*[*rule-format*]:

*sorted-spmat*  $A \longrightarrow$  *sorted-spvec*  $B \longrightarrow$  *sparse-row-matrix* (*mult-spmat*  $A B$ ) =  
(*sparse-row-matrix*  $A$ ) \* (*sparse-row-matrix*  $B$ )  
⟨*proof*⟩

**lemma** *sorted-spvec-mult-spmat*[*rule-format*]:  
*sorted-spvec* ( $A::('a::\text{ordered-ring}) \text{ spmat}$ )  $\longrightarrow$  *sorted-spvec* (*mult-spmat*  $A B$ )  
⟨*proof*⟩

**lemma** *sorted-spmat-mult-spmat*[*rule-format*]:  
*sorted-spmat* ( $B::('a::\text{ordered-ring}) \text{ spmat}$ )  $\longrightarrow$  *sorted-spmat* (*mult-spmat*  $A B$ )  
⟨*proof*⟩

**consts**

*add-spvec* :: ( $'a::\text{ordered-ab-group}$ ) *spvec* \*  $'a \text{ spvec} \Rightarrow 'a \text{ spvec}$   
*add-spmat* :: ( $'a::\text{ordered-ab-group}$ ) *spmat* \*  $'a \text{ spmat} \Rightarrow 'a \text{ spmat}$

**recdef** *add-spvec measure* (% ( $a, b$ ). *length*  $a$  + (*length*  $b$ ))  
*add-spvec* ( $arr, []$ ) =  $arr$   
*add-spvec* ( $[], brr$ ) =  $brr$   
*add-spvec* ( $a\#arr, b\#brr$ ) = (  
*if* (*fst*  $a$ ) < (*fst*  $b$ ) *then* ( $a\#(\text{add-spvec } (arr, b\#brr))$ )  
*else* (*if* (*fst*  $b$ ) < (*fst*  $a$ ) *then* ( $b\#(\text{add-spvec } (a\#arr, brr))$ )  
*else* ( $(\text{fst } a, (\text{snd } a) + (\text{snd } b))\#(\text{add-spvec } (arr, brr))$ ))

**lemma** *add-spvec-empty1*[*simp*]: *add-spvec* ( $[], a$ ) =  $a$   
⟨*proof*⟩

**lemma** *add-spvec-empty2*[*simp*]: *add-spvec* ( $a, []$ ) =  $a$   
⟨*proof*⟩

**lemma** *sparse-row-vector-add*: *sparse-row-vector* (*add-spvec* ( $a, b$ )) = (*sparse-row-vector*  $a$ ) + (*sparse-row-vector*  $b$ )  
⟨*proof*⟩

**recdef** *add-spmat measure* (% ( $A, B$ ). (*length*  $A$ ) + (*length*  $B$ ))  
*add-spmat* ( $[], bs$ ) =  $bs$   
*add-spmat* ( $as, []$ ) =  $as$   
*add-spmat* ( $a\#as, b\#bs$ ) = (  
*if* *fst*  $a$  < *fst*  $b$  *then*  
( $a\#(\text{add-spmat } (as, b\#bs))$ )  
*else* (*if* *fst*  $b$  < *fst*  $a$  *then*  
( $b\#(\text{add-spmat } (a\#as, bs))$ )  
*else*  
( $(\text{fst } a, \text{add-spvec } (\text{snd } a, \text{snd } b))\#(\text{add-spmat } (as, bs))$ ))

**lemma** *sparse-row-add-spmat*: *sparse-row-matrix* (*add-spmat* ( $A, B$ )) = (*sparse-row-matrix*  $A$ ) + (*sparse-row-matrix*  $B$ )  
⟨*proof*⟩

**lemma** *sorted-add-spvec-helper1*[rule-format]:  $add\_spvec ((a,b)\#arr, brr) = (ab, bb) \# list \longrightarrow (ab = a \mid (brr \neq [] \ \& \ ab = fst (hd \ brr)))$   
 ⟨proof⟩

**lemma** *sorted-add-spmat-helper1*[rule-format]:  $add\_spmat ((a,b)\#arr, brr) = (ab, bb) \# list \longrightarrow (ab = a \mid (brr \neq [] \ \& \ ab = fst (hd \ brr)))$   
 ⟨proof⟩

**lemma** *sorted-add-spvec-helper*[rule-format]:  $add\_spvec (arr, brr) = (ab, bb) \# list \longrightarrow ((arr \neq [] \ \& \ ab = fst (hd \ arr)) \mid (brr \neq [] \ \& \ ab = fst (hd \ brr)))$   
 ⟨proof⟩

**lemma** *sorted-add-spmat-helper*[rule-format]:  $add\_spmat (arr, brr) = (ab, bb) \# list \longrightarrow ((arr \neq [] \ \& \ ab = fst (hd \ arr)) \mid (brr \neq [] \ \& \ ab = fst (hd \ brr)))$   
 ⟨proof⟩

**lemma** *add-spvec-commute*:  $add\_spvec (a, b) = add\_spvec (b, a)$   
 ⟨proof⟩

**lemma** *add-spmat-commute*:  $add\_spmat (a, b) = add\_spmat (b, a)$   
 ⟨proof⟩

**lemma** *sorted-add-spvec-helper2*:  $add\_spvec ((a,b)\#arr, brr) = (ab, bb) \# list \implies aa < a \implies sorted\_spvec ((aa, ba) \# brr) \implies aa < ab$   
 ⟨proof⟩

**lemma** *sorted-add-spmat-helper2*:  $add\_spmat ((a,b)\#arr, brr) = (ab, bb) \# list \implies aa < a \implies sorted\_spvec ((aa, ba) \# brr) \implies aa < ab$   
 ⟨proof⟩

**lemma** *sorted-spvec-add-spvec*[rule-format]:  $sorted\_spvec \ a \longrightarrow sorted\_spvec \ b \longrightarrow sorted\_spvec (add\_spvec (a, b))$   
 ⟨proof⟩

**lemma** *sorted-spvec-add-spmat*[rule-format]:  $sorted\_spvec \ A \longrightarrow sorted\_spvec \ B \longrightarrow sorted\_spvec (add\_spmat (A, B))$   
 ⟨proof⟩

**lemma** *sorted-spmat-add-spmat*[rule-format]:  $sorted\_spmat \ A \longrightarrow sorted\_spmat \ B \longrightarrow sorted\_spmat (add\_spmat (A, B))$   
 ⟨proof⟩

**consts**

$le\_spvec :: ('a::ordered-ab-group) \ spvec * 'a \ spvec \Rightarrow bool$   
 $le\_spmat :: ('a::ordered-ab-group) \ spmat * 'a \ spmat \Rightarrow bool$

**recdef** *le-spvec measure* (% (a,b). (length a) + (length b))  
 $le\_spvec ([] , []) = True$   
 $le\_spvec (a\#as, []) = ((snd \ a \leq 0) \ \& \ (le\_spvec (as, [])))$

```

le-spvec ([], b#bs) = ((0 <= snd b) & (le-spvec ([], bs)))
le-spvec (a#as, b#bs) = (
  if (fst a < fst b) then
    ((snd a <= 0) & (le-spvec (as, b#bs)))
  else (if (fst b < fst a) then
    ((0 <= snd b) & (le-spvec (a#as, bs)))
  else
    ((snd a <= snd b) & (le-spvec (as, bs))))

```

```

recdef le-spmat measure (% (a,b). (length a) + (length b))
le-spmat ([], []) = True
le-spmat (a#as, []) = (le-spvec (snd a, []) & (le-spmat (as, [])))
le-spmat ([], b#bs) = (le-spvec ([], snd b) & (le-spmat ([], bs)))
le-spmat (a#as, b#bs) = (
  if fst a < fst b then
    (le-spvec(snd a,[]) & le-spmat(as, b#bs))
  else (if (fst b < fst a) then
    (le-spvec([], snd b) & le-spmat(a#as, bs))
  else
    (le-spvec(snd a, snd b) & le-spmat (as, bs)))

```

### constdefs

```

disj-matrices :: ('a::zero) matrix => 'a matrix => bool
disj-matrices A B == (! j i. (Rep-matrix A j i ≠ 0) → (Rep-matrix B j i = 0)) & (! j i. (Rep-matrix B j i ≠ 0) → (Rep-matrix A j i = 0))

```

⟨ML⟩

```

lemma disj-matrices-contr1: disj-matrices A B ⇒ Rep-matrix A j i ≠ 0 ⇒ Rep-matrix B j i = 0
  ⟨proof⟩

```

```

lemma disj-matrices-contr2: disj-matrices A B ⇒ Rep-matrix B j i ≠ 0 ⇒ Rep-matrix A j i = 0
  ⟨proof⟩

```

```

lemma disj-matrices-add: disj-matrices A B ⇒ disj-matrices C D ⇒ disj-matrices A D ⇒ disj-matrices B C ⇒
  (A + B <= C + D) = (A <= C & B <= (D::('a::lordered-ab-group) matrix))
  ⟨proof⟩

```

```

lemma disj-matrices-zero1[simp]: disj-matrices 0 B
  ⟨proof⟩

```

```

lemma disj-matrices-zero2[simp]: disj-matrices A 0
  ⟨proof⟩

```

```

lemma disj-matrices-commute: disj-matrices A B = disj-matrices B A

```

*<proof>*

**lemma** *disj-matrices-add-le-zero*:  $\text{disj-matrices } A \ B \implies$   
 $(A + B \leq 0) = (A \leq 0 \ \& \ (B::('a::\text{lordered-ab-group}) \text{matrix}) \leq 0)$   
*<proof>*

**lemma** *disj-matrices-add-zero-le*:  $\text{disj-matrices } A \ B \implies$   
 $(0 \leq A + B) = (0 \leq A \ \& \ 0 \leq (B::('a::\text{lordered-ab-group}) \text{matrix}))$   
*<proof>*

**lemma** *disj-matrices-add-x-le*:  $\text{disj-matrices } A \ B \implies \text{disj-matrices } B \ C \implies$   
 $(A \leq B + C) = (A \leq C \ \& \ 0 \leq (B::('a::\text{lordered-ab-group}) \text{matrix}))$   
*<proof>*

**lemma** *disj-matrices-add-le-x*:  $\text{disj-matrices } A \ B \implies \text{disj-matrices } B \ C \implies$   
 $(B + A \leq C) = (A \leq C \ \& \ (B::('a::\text{lordered-ab-group}) \text{matrix}) \leq 0)$   
*<proof>*

**lemma** *disj-sparse-row-singleton*:  $i \leq j \implies \text{sorted-spvec}((j,y)\#v) \implies \text{disj-matrices}$   
 $(\text{sparse-row-vector } v) \ (\text{singleton-matrix } 0 \ i \ x)$   
*<proof>*

**lemma** *disj-matrices-x-add*:  $\text{disj-matrices } A \ B \implies \text{disj-matrices } A \ C \implies \text{disj-matrices}$   
 $(A::('a::\text{lordered-ab-group}) \text{matrix}) \ (B+C)$   
*<proof>*

**lemma** *disj-matrices-add-x*:  $\text{disj-matrices } A \ B \implies \text{disj-matrices } A \ C \implies \text{disj-matrices}$   
 $(B+C) \ (A::('a::\text{lordered-ab-group}) \text{matrix})$   
*<proof>*

**lemma** *disj-singleton-matrices[simp]*:  $\text{disj-matrices} \ (\text{singleton-matrix } j \ i \ x) \ (\text{singleton-matrix}$   
 $u \ v \ y) = (j \neq u \ | \ i \neq v \ | \ x = 0 \ | \ y = 0)$   
*<proof>*

**lemma** *disj-move-sparse-vec-mat[simplified disj-matrices-commute]*:  
 $j \leq a \implies \text{sorted-spvec}((a,c)\#as) \implies \text{disj-matrices} \ (\text{move-matrix} \ (\text{sparse-row-vector}$   
 $b) \ (\text{int } j) \ i) \ (\text{sparse-row-matrix } as)$   
*<proof>*

**lemma** *disj-move-sparse-row-vector-twice*:  
 $j \neq u \implies \text{disj-matrices} \ (\text{move-matrix} \ (\text{sparse-row-vector } a) \ j \ i) \ (\text{move-matrix}$   
 $(\text{sparse-row-vector } b) \ u \ v)$   
*<proof>*

**lemma** *le-spvec-iff-sparse-row-le[rule-format]*:  $(\text{sorted-spvec } a) \longrightarrow (\text{sorted-spvec}$   
 $b) \longrightarrow (\text{le-spvec } (a,b)) = (\text{sparse-row-vector } a \leq \text{sparse-row-vector } b)$   
*<proof>*

**lemma** *le-spvec-empty2-sparse-row[rule-format]*:  $(\text{sorted-spvec } b) \longrightarrow (\text{le-spvec } (b,[]))$

= (*sparse-row-vector*  $b \leq 0$ )  
 ⟨*proof*⟩

**lemma** *le-spvec-empty1-sparse-row*[*rule-format*]: (*sorted-spvec*  $b$ )  $\longrightarrow$  (*le-spvec* ( $\square, b$ )  
 = ( $0 \leq$  *sparse-row-vector*  $b$ )  
 ⟨*proof*⟩

**lemma** *le-spmat-iff-sparse-row-le*[*rule-format*]: (*sorted-spvec*  $A$ )  $\longrightarrow$  (*sorted-spmat*  
 $A$ )  $\longrightarrow$  (*sorted-spvec*  $B$ )  $\longrightarrow$  (*sorted-spmat*  $B$ )  $\longrightarrow$   
*le-spmat*( $A, B$ ) = (*sparse-row-matrix*  $A \leq$  *sparse-row-matrix*  $B$ )  
 ⟨*proof*⟩

⟨*ML*⟩

**consts**

*abs-spmat* :: ('*a*::*lordered-ring*) *spmat*  $\Rightarrow$  '*a* *spmat*  
*minus-spmat* :: ('*a*::*lordered-ring*) *spmat*  $\Rightarrow$  '*a* *spmat*

**primrec**

*abs-spmat*  $\square$  =  $\square$   
*abs-spmat* ( $a\#as$ ) = (*fst*  $a$ , *abs-spvec* (*snd*  $a$ ))#(*abs-spmat*  $as$ )

**primrec**

*minus-spmat*  $\square$  =  $\square$   
*minus-spmat* ( $a\#as$ ) = (*fst*  $a$ , *minus-spvec* (*snd*  $a$ ))#(*minus-spmat*  $as$ )

**lemma** *sparse-row-matrix-minus*:

*sparse-row-matrix* (*minus-spmat*  $A$ ) = - (*sparse-row-matrix*  $A$ )  
 ⟨*proof*⟩

**lemma** *Rep-sparse-row-vector-zero*:  $x \neq 0 \implies$  *Rep-matrix* (*sparse-row-vector*  $v$ )  
 $x y = 0$   
 ⟨*proof*⟩

**lemma** *sparse-row-matrix-abs*:

*sorted-spvec*  $A \implies$  *sorted-spmat*  $A \implies$  *sparse-row-matrix* (*abs-spmat*  $A$ ) = *abs*  
 (*sparse-row-matrix*  $A$ )  
 ⟨*proof*⟩

**lemma** *sorted-spvec-minus-spmat*: *sorted-spvec*  $A \implies$  *sorted-spvec* (*minus-spmat*  
 $A$ )  
 ⟨*proof*⟩

**lemma** *sorted-spvec-abs-spmat*: *sorted-spvec*  $A \implies$  *sorted-spvec* (*abs-spmat*  $A$ )  
 ⟨*proof*⟩

**lemma** *sorted-spmat-minus-spmat*: *sorted-spmat*  $A \implies$  *sorted-spmat* (*minus-spmat*  
 $A$ )  
 ⟨*proof*⟩

**lemma** *sorted-spmat-abs-spmat*:  $\text{sorted-spmat } A \implies \text{sorted-spmat } (\text{abs-spmat } A)$   
*<proof>*

**constdefs**

*diff-spmat* :: ('a::lordered-ring) *spmat*  $\Rightarrow$  'a *spmat*  $\Rightarrow$  'a *spmat*  
*diff-spmat* A B == *add-spmat* (A, *minus-spmat* B)

**lemma** *sorted-spmat-diff-spmat*:  $\text{sorted-spmat } A \implies \text{sorted-spmat } B \implies \text{sorted-spmat } (\text{diff-spmat } A B)$   
*<proof>*

**lemma** *sorted-spmat-diff-spmat*:  $\text{sorted-spmat } A \implies \text{sorted-spmat } B \implies \text{sorted-spmat } (\text{diff-spmat } A B)$   
*<proof>*

**lemma** *sparse-row-diff-spmat*:  $\text{sparse-row-matrix } (\text{diff-spmat } A B) = (\text{sparse-row-matrix } A) - (\text{sparse-row-matrix } B)$   
*<proof>*

**constdefs**

*sorted-sparse-matrix* :: 'a *spmat*  $\Rightarrow$  bool  
*sorted-sparse-matrix* A == (*sorted-spmat* A) & (*sorted-spmat* A)

**lemma** *sorted-sparse-matrix-imp-spmat*:  $\text{sorted-sparse-matrix } A \implies \text{sorted-spmat } A$   
*<proof>*

**lemma** *sorted-sparse-matrix-imp-spmat*:  $\text{sorted-sparse-matrix } A \implies \text{sorted-spmat } A$   
*<proof>*

**lemmas** *sorted-sp-simps* =  
*sorted-spmat.simps*  
*sorted-spmat.simps*  
*sorted-sparse-matrix-def*

**lemma** *bool1*:  $(\neg \text{True}) = \text{False}$  *<proof>*

**lemma** *bool2*:  $(\neg \text{False}) = \text{True}$  *<proof>*

**lemma** *bool3*:  $((P::\text{bool}) \wedge \text{True}) = P$  *<proof>*

**lemma** *bool4*:  $(\text{True} \wedge (P::\text{bool})) = P$  *<proof>*

**lemma** *bool5*:  $((P::\text{bool}) \wedge \text{False}) = \text{False}$  *<proof>*

**lemma** *bool6*:  $(\text{False} \wedge (P::\text{bool})) = \text{False}$  *<proof>*

**lemma** *bool7*:  $((P::\text{bool}) \vee \text{True}) = \text{True}$  *<proof>*

**lemma** *bool8*:  $(\text{True} \vee (P::\text{bool})) = \text{True}$  *<proof>*

**lemma** *bool9*:  $((P::\text{bool}) \vee \text{False}) = P$  *<proof>*

**lemma** *bool10*:  $(\text{False} \vee (P::\text{bool})) = P$  *<proof>*

**lemmas** *boolarith* = *bool1 bool2 bool3 bool4 bool5 bool6 bool7 bool8 bool9 bool10*

**lemma** *if-case-eq*:  $(\text{if } b \text{ then } x \text{ else } y) = (\text{case } b \text{ of } \text{True} \Rightarrow x \mid \text{False} \Rightarrow y)$

*<proof>*

**consts**

$pprt\text{-}spvec :: ('a::\{lordered\text{-}ab\text{-}group\})\ spvec \Rightarrow 'a\ spvec$   
 $np\text{r}t\text{-}spvec :: ('a::\{lordered\text{-}ab\text{-}group\})\ spvec \Rightarrow 'a\ spvec$   
 $pprt\text{-}spmat :: ('a::\{lordered\text{-}ab\text{-}group\})\ smat \Rightarrow 'a\ smat$   
 $np\text{r}t\text{-}spmat :: ('a::\{lordered\text{-}ab\text{-}group\})\ smat \Rightarrow 'a\ smat$

**primrec**

$pprt\text{-}spvec\ [] = []$   
 $pprt\text{-}spvec\ (a\#\ as) = (fst\ a,\ pprt\ (snd\ a))\ \#\ (pprt\text{-}spvec\ as)$

**primrec**

$np\text{r}t\text{-}spvec\ [] = []$   
 $np\text{r}t\text{-}spvec\ (a\#\ as) = (fst\ a,\ np\text{r}t\ (snd\ a))\ \#\ (np\text{r}t\text{-}spvec\ as)$

**primrec**

$pprt\text{-}spmat\ [] = []$   
 $pprt\text{-}spmat\ (a\#\ as) = (fst\ a,\ pprt\text{-}spvec\ (snd\ a))\ \#\ (pprt\text{-}spmat\ as)$

**primrec**

$np\text{r}t\text{-}spmat\ [] = []$   
 $np\text{r}t\text{-}spmat\ (a\#\ as) = (fst\ a,\ np\text{r}t\text{-}spvec\ (snd\ a))\ \#\ (np\text{r}t\text{-}spmat\ as)$

**lemma** *pprt-add: disj-matrices*  $A\ (B::(-::lordered\text{-}ring)\ matrix) \Longrightarrow pprt\ (A+B)$   
 $= pprt\ A + pprt\ B$   
*<proof>*

**lemma** *np\text{r}t-add: disj-matrices*  $A\ (B::(-::lordered\text{-}ring)\ matrix) \Longrightarrow np\text{r}t\ (A+B)$   
 $= np\text{r}t\ A + np\text{r}t\ B$   
*<proof>*

**lemma** *pprt-singleton[simp]:*  $pprt\ (singleton\text{-}matrix\ j\ i\ (x::(-::lordered\text{-}ring))) = singleton\text{-}matrix$   
 $j\ i\ (pprt\ x)$   
*<proof>*

**lemma** *np\text{r}t-singleton[simp]:*  $np\text{r}t\ (singleton\text{-}matrix\ j\ i\ (x::(-::lordered\text{-}ring))) = singleton\text{-}matrix$   
 $j\ i\ (np\text{r}t\ x)$   
*<proof>*

**lemma** *less-imp-le:*  $a < b \Longrightarrow a \leq (b::(-::order))$  *<proof>*

**lemma** *sparse-row-vector-pprt:*  $sorted\text{-}spvec\ v \Longrightarrow sparse\text{-}row\text{-}vector\ (pprt\text{-}spvec$   
 $v) = pprt\ (sparse\text{-}row\text{-}vector\ v)$   
*<proof>*

**lemma** *sparse-row-vector-nprt*:  $\text{sorted-spvec } v \implies \text{sparse-row-vector } (\text{nprt-spvec } v) = \text{nprt } (\text{sparse-row-vector } v)$   
 ⟨proof⟩

**lemma** *pprt-move-matrix*:  $\text{pprt } (\text{move-matrix } (A::('a::lordered-ring) \text{ matrix}) j i) = \text{move-matrix } (\text{pprt } A) j i$   
 ⟨proof⟩

**lemma** *nprt-move-matrix*:  $\text{nprt } (\text{move-matrix } (A::('a::lordered-ring) \text{ matrix}) j i) = \text{move-matrix } (\text{nprt } A) j i$   
 ⟨proof⟩

**lemma** *sparse-row-matrix-pprt*:  $\text{sorted-spvec } m \implies \text{sorted-spmat } m \implies \text{sparse-row-matrix } (\text{pprt-spmat } m) = \text{pprt } (\text{sparse-row-matrix } m)$   
 ⟨proof⟩

**lemma** *sparse-row-matrix-nprt*:  $\text{sorted-spvec } m \implies \text{sorted-spmat } m \implies \text{sparse-row-matrix } (\text{nprt-spmat } m) = \text{nprt } (\text{sparse-row-matrix } m)$   
 ⟨proof⟩

**lemma** *sorted-pprt-spvec*:  $\text{sorted-spvec } v \implies \text{sorted-spvec } (\text{pprt-spvec } v)$   
 ⟨proof⟩

**lemma** *sorted-nprt-spvec*:  $\text{sorted-spvec } v \implies \text{sorted-spvec } (\text{nprt-spvec } v)$   
 ⟨proof⟩

**lemma** *sorted-spvec-pprt-spmat*:  $\text{sorted-spvec } m \implies \text{sorted-spvec } (\text{pprt-spmat } m)$   
 ⟨proof⟩

**lemma** *sorted-spvec-nprt-spmat*:  $\text{sorted-spvec } m \implies \text{sorted-spvec } (\text{nprt-spmat } m)$   
 ⟨proof⟩

**lemma** *sorted-spmat-pprt-spmat*:  $\text{sorted-spmat } m \implies \text{sorted-spmat } (\text{pprt-spmat } m)$   
 ⟨proof⟩

**lemma** *sorted-spmat-nprt-spmat*:  $\text{sorted-spmat } m \implies \text{sorted-spmat } (\text{nprt-spmat } m)$   
 ⟨proof⟩

**constdefs**

$\text{mult-est-spmat} :: ('a::lordered-ring) \text{ spmat} \Rightarrow 'a \text{ spmat} \Rightarrow 'a \text{ spmat} \Rightarrow 'a \text{ spmat} \Rightarrow 'a \text{ spmat}$

$\text{mult-est-spmat } r1 \ r2 \ s1 \ s2 ==$   
 $\text{add-spmat } (\text{mult-spmat } (\text{pprt-spmat } s2) (\text{pprt-spmat } r2), \text{add-spmat } (\text{mult-spmat } (\text{pprt-spmat } s1) (\text{nprt-spmat } r2),$   
 $\text{add-spmat } (\text{mult-spmat } (\text{nprt-spmat } s2) (\text{pprt-spmat } r1), \text{mult-spmat } (\text{nprt-spmat } s1) (\text{nprt-spmat } r1))))$

**lemmas** *sparse-row-matrix-op-simps* =  
*sorted-sparse-matrix-imp-spmat sorted-sparse-matrix-imp-spvec*  
*sparse-row-add-spmat sorted-spvec-add-spmat sorted-spmat-add-spmat*  
*sparse-row-diff-spmat sorted-spvec-diff-spmat sorted-spmat-diff-spmat*  
*sparse-row-matrix-minus sorted-spvec-minus-spmat sorted-spmat-minus-spmat*  
*sparse-row-mult-spmat sorted-spvec-mult-spmat sorted-spmat-mult-spmat*  
*sparse-row-matrix-abs sorted-spvec-abs-spmat sorted-spmat-abs-spmat*  
*le-spmat-iff-sparse-row-le*  
*sparse-row-matrix-pprt sorted-spvec-pprt-spmat sorted-spmat-pprt-spmat*  
*sparse-row-matrix-nprt sorted-spvec-nprt-spmat sorted-spmat-nprt-spmat*

**lemma** *zero-eq-Numerals*:  $(0::\text{number-ring}) = \text{Numerals } 0$  *<proof>*

**lemmas** *sparse-row-matrix-arith-simps*[*simplified zero-eq-Numerals*] =  
*mult-spmat.simps mult-spvec-spmat.simps*  
*addmult-spvec.simps*  
*smult-spvec-empty smult-spvec-cons*  
*add-spmat.simps add-spvec.simps*  
*minus-spmat.simps minus-spvec.simps*  
*abs-spmat.simps abs-spvec.simps*  
*diff-spmat-def*  
*le-spmat.simps le-spvec.simps*  
*pprt-spmat.simps pprt-spvec.simps*  
*nprt-spmat.simps nprt-spvec.simps*  
*mult-est-spmat-def*

**lemma** *spm-mult-le-dual-prts*:

**assumes**

*sorted-sparse-matrix A1*  
*sorted-sparse-matrix A2*  
*sorted-sparse-matrix c1*  
*sorted-sparse-matrix c2*  
*sorted-sparse-matrix y*  
*sorted-sparse-matrix r1*  
*sorted-sparse-matrix r2*  
*sorted-spvec b*  
*le-spmat ([], y)*  
*sparse-row-matrix A1  $\leq$  A*  
*A  $\leq$  sparse-row-matrix A2*  
*sparse-row-matrix c1  $\leq$  c*  
*c  $\leq$  sparse-row-matrix c2*  
*sparse-row-matrix r1  $\leq$  x*  
*x  $\leq$  sparse-row-matrix r2*  
*A \* x  $\leq$  sparse-row-matrix (b::('a::lordered-ring) spmat)*

**shows**

```

  c * x ≤ sparse-row-matrix (add-spmat (mult-spmat y b,
    (let s1 = diff-spmat c1 (mult-spmat y A2); s2 = diff-spmat c2 (mult-spmat y
A1) in
      add-spmat (mult-spmat (pprt-spmat s2) (pprt-spmat r2), add-spmat (mult-spmat
(pprt-spmat s1) (nprrt-spmat r2),
      add-spmat (mult-spmat (nprrt-spmat s2) (pprt-spmat r1), mult-spmat (nprrt-spmat
s1) (nprrt-spmat r1))))))
  ⟨proof⟩

```

**lemma** *spm-mult-le-dual-prts-no-let*:

```

assumes
  sorted-sparse-matrix A1
  sorted-sparse-matrix A2
  sorted-sparse-matrix c1
  sorted-sparse-matrix c2
  sorted-sparse-matrix y
  sorted-sparse-matrix r1
  sorted-sparse-matrix r2
  sorted-spvec b
  le-spmat ([], y)
  sparse-row-matrix A1 ≤ A
  A ≤ sparse-row-matrix A2
  sparse-row-matrix c1 ≤ c
  c ≤ sparse-row-matrix c2
  sparse-row-matrix r1 ≤ x
  x ≤ sparse-row-matrix r2
  A * x ≤ sparse-row-matrix (b::('a::lordered-ring) spmat)
shows
  c * x ≤ sparse-row-matrix (add-spmat (mult-spmat y b,
  mult-est-spmat r1 r2 (diff-spmat c1 (mult-spmat y A2)) (diff-spmat c2 (mult-spmat
y A1))))
  ⟨proof⟩

```

**end**

**theory** *FloatSparseMatrix* **imports** *Float SparseMatrix* **begin**

**end**

```

theory Cplex
imports FloatSparseMatrix
uses Cplex-tools.ML CplexMatrixConverter.ML FloatSparseMatrixBuilder.ML fspmlp.ML
begin

```

**end**

```

theory MatrixLP
imports Cplex
begin

```

```

constdefs

```

```

  list-case-compute :: 'b list ⇒ 'a ⇒ ('b ⇒ 'b list ⇒ 'a) ⇒ 'a
  list-case-compute l a f == list-case a f l

```

```

lemma list-case-compute: list-case = (λ (a::'a) f (l::'b list). list-case-compute l a
f)
  ⟨proof⟩

```

```

lemma list-case-compute-empty: list-case-compute ([]::'b list) = (λ (a::'a) f. a)
  ⟨proof⟩

```

```

lemma list-case-compute-cons: list-case-compute (u#v) = (λ (a::'a) f. (f (u::'b)
v))
  ⟨proof⟩

```

```

lemma If-True: (If True) = (λ x y. x)
  ⟨proof⟩

```

```

lemma If-False: (If False) = (λ x y. y)
  ⟨proof⟩

```

```

lemma Let-compute: Let (x::'a) f = ((f x)::'b)
  ⟨proof⟩

```

```

lemma fst-compute: fst (a::'a, b::'b) = a
  ⟨proof⟩

```

```

lemma snd-compute: snd (a::'a, b::'b) = b
  ⟨proof⟩

```

```

lemma bool1: (¬ True) = False ⟨proof⟩

```

```

lemma bool2: (¬ False) = True ⟨proof⟩

```

```

lemma bool3: ((P::bool) ∧ True) = P ⟨proof⟩

```

```

lemma bool4: (True ∧ (P::bool)) = P ⟨proof⟩

```

```

lemma bool5: ((P::bool) ∧ False) = False ⟨proof⟩

```

```

lemma bool6: (False ∧ (P::bool)) = False ⟨proof⟩

```

```

lemma bool7: ((P::bool) ∨ True) = True ⟨proof⟩

```

```

lemma bool8: (True ∨ (P::bool)) = True ⟨proof⟩

```

```

lemma bool9: ((P::bool) ∨ False) = P ⟨proof⟩

```

```
lemma bool10: (False  $\vee$  (P::bool)) = P <proof>  
lemmas boolarith = bool1 bool2 bool3 bool4 bool5 bool6 bool7 bool8 bool9 bool10  
  
lemmas float-arith = Float.arith  
lemmas sparse-row-matrix-arith-simps = SparseMatrix.sparse-row-matrix-arith-simps  
lemmas sorted-sp-simps = SparseMatrix.sorted-sp-simps  
lemmas fst-snd-conv = Product-Type.fst-conv Product-Type.snd-conv  
  
end
```