

Isabelle/FOL — First-Order Logic

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Contents

1	Intuitionistic first-order logic	1
1.1	Syntax and axiomatic basis	2
1.2	Lemmas and proof tools	4
1.3	Intuitionistic Reasoning	11
1.4	Atomizing meta-level rules	11
1.5	Atomizing elimination rules	12
1.6	Calculational rules	12
1.7	“Let” declarations	12
1.8	Intuitionistic simplification rules	13
1.9	Legacy ML bindings	15
2	Classical first-order logic	15
2.1	The classical axiom	15
2.2	Lemmas and proof tools	16
2.3	Other simple lemmas	20
2.4	Proof by cases and induction	20

1 Intuitionistic first-order logic

```
theory IFOL
imports Pure
uses
  ~~ /src/Provers/splitter.ML
  ~~ /src/Provers/hypsubst.ML
  ~~ /src/Tools/IsaPlanner/zipper.ML
  ~~ /src/Tools/IsaPlanner/isand.ML
  ~~ /src/Tools/IsaPlanner/rw-tools.ML
  ~~ /src/Tools/IsaPlanner/rw-inst.ML
  ~~ /src/Provers/eqsubst.ML
  ~~ /src/Provers/quantifier1.ML
  ~~ /src/Provers/project-rule.ML
  ~~ /src/Tools/atomize-elim.ML
```

```

(fologic.ML)
(hypsubstdata.ML)
(intprover.ML)
begin

```

1.1 Syntax and axiomatic basis

$\langle ML \rangle$

global

classes *term*
defaultsort *term*

typeddecl *o*

judgment
Trueprop $:: o \Rightarrow prop$ $((-) 5)$

consts
True $:: o$
False $:: o$

op = $:: ['a, 'a] \Rightarrow o$ **(infixl = 50)**

Not $:: o \Rightarrow o$ $(\sim - [40] 40)$
op & $:: [o, o] \Rightarrow o$ **(infixr & 35)**
op | $:: [o, o] \Rightarrow o$ **(infixr | 30)**
op $-->$ $:: [o, o] \Rightarrow o$ **(infixr $-->$ 25)**
op $<->$ $:: [o, o] \Rightarrow o$ **(infixr $<->$ 25)**

All $:: ('a \Rightarrow o) \Rightarrow o$ **(binder ALL 10)**
Ex $:: ('a \Rightarrow o) \Rightarrow o$ **(binder EX 10)**
Ex1 $:: ('a \Rightarrow o) \Rightarrow o$ **(binder EX! 10)**

abbreviation

not-equal $:: ['a, 'a] \Rightarrow o$ **(infixl $\sim =$ 50)** **where**
 $x \sim = y == \sim (x = y)$

notation (*xsymbols*)

not-equal **(infixl \neq 50)**

notation (*HTML output*)

not-equal **(infixl \neq 50)**

notation (*xsymbols*)
Not (\neg - [40] 40) and
op & (**infixr** \wedge 35) and
op | (**infixr** \vee 30) and
All (**binder** \forall 10) and
Ex (**binder** \exists 10) and
Ex1 (**binder** $\exists!$ 10) and
op $-->$ (**infixr** \longrightarrow 25) and
op $<->$ (**infixr** \longleftrightarrow 25)

notation (*HTML output*)
Not (\neg - [40] 40) and
op & (**infixr** \wedge 35) and
op | (**infixr** \vee 30) and
All (**binder** \forall 10) and
Ex (**binder** \exists 10) and
Ex1 (**binder** $\exists!$ 10)

local

finalconsts

False All Ex
op =
op &
op |
op $-->$

axioms

refl: $a=a$

conjI: $[| P; Q |] ==> P \& Q$
conjunct1: $P \& Q ==> P$
conjunct2: $P \& Q ==> Q$

disjI1: $P ==> P | Q$
disjI2: $Q ==> P | Q$
disjE: $[| P | Q; P ==> R; Q ==> R |] ==> R$

impI: $(P ==> Q) ==> P --> Q$
mp: $[| P --> Q; P |] ==> Q$

FalseE: $False ==> P$

$allI: \quad (!x. P(x)) \implies (ALL\ x. P(x))$
 $spec: \quad (ALL\ x. P(x)) \implies P(x)$

 $exI: \quad P(x) \implies (EX\ x. P(x))$
 $exE: \quad [| EX\ x. P(x); !x. P(x) \implies R |] \implies R$

$eq\text{-}reflection: \quad (x=y) \implies (x==y)$
 $iff\text{-}reflection: \quad (P<->Q) \implies (P==Q)$

lemmas *strip* = *impI allI*

Thanks to Stephan Merz

theorem *subst*:
 assumes *eq*: $a = b$ **and** $p: P(a)$
 shows $P(b)$
 $\langle proof \rangle$

defs

$True\text{-}def: \quad True == False \dashv\dashv False$
 $not\text{-}def: \quad \sim P == P \dashv\dashv False$
 $iff\text{-}def: \quad P<->Q == (P \dashv\dashv Q) \ \& \ (Q \dashv\dashv P)$

$ex1\text{-}def: \quad Ex1(P) == EX\ x. P(x) \ \& \ (ALL\ y. P(y) \dashv\dashv y=x)$

1.2 Lemmas and proof tools

lemma *TrueI*: *True*
 $\langle proof \rangle$

lemma *conjE*:
 assumes *major*: $P \ \& \ Q$
 and $r: [| P; Q |] \implies R$
 shows R
 $\langle proof \rangle$

lemma *impE*:

assumes *major*: $P \multimap Q$
and P
and r : $Q \implies R$
shows R
 $\langle proof \rangle$

lemma *allE*:
assumes *major*: $\text{ALL } x. P(x)$
and r : $P(x) \implies R$
shows R
 $\langle proof \rangle$

lemma *all-dupE*:
assumes *major*: $\text{ALL } x. P(x)$
and r : $[| P(x); \text{ALL } x. P(x) |] \implies R$
shows R
 $\langle proof \rangle$

lemma *notI*: $(P \implies \text{False}) \implies \sim P$
 $\langle proof \rangle$

lemma *notE*: $[| \sim P; P |] \implies R$
 $\langle proof \rangle$

lemma *rev-notE*: $[| P; \sim P |] \implies R$
 $\langle proof \rangle$

lemma *not-to-imp*:
assumes $\sim P$
and r : $P \multimap \text{False} \implies Q$
shows Q
 $\langle proof \rangle$

lemma *rev-mp*: $[| P; P \multimap Q |] \implies Q$
 $\langle proof \rangle$

lemma *contrapos*:
assumes *major*: $\sim Q$
and *minor*: $P \implies Q$
shows $\sim P$
 $\langle proof \rangle$

$\langle ML \rangle$

lemma *iffI*: $[| P \implies Q; Q \implies P |] \implies P \leftrightarrow Q$
 $\langle proof \rangle$

lemma *iffE*:
 assumes *major*: $P \leftrightarrow Q$
 and *r*: $P \rightarrow Q \implies Q \rightarrow P \implies R$
 shows *R*
 $\langle proof \rangle$

lemma *iffD1*: $[| P \leftrightarrow Q; P |] \implies Q$
 $\langle proof \rangle$

lemma *iffD2*: $[| P \leftrightarrow Q; Q |] \implies P$
 $\langle proof \rangle$

lemma *rev-iffD1*: $[| P; P \leftrightarrow Q |] \implies Q$
 $\langle proof \rangle$

lemma *rev-iffD2*: $[| Q; P \leftrightarrow Q |] \implies P$
 $\langle proof \rangle$

lemma *iff-refl*: $P \leftrightarrow P$
 $\langle proof \rangle$

lemma *iff-sym*: $Q \leftrightarrow P \implies P \leftrightarrow Q$
 $\langle proof \rangle$

lemma *iff-trans*: $[| P \leftrightarrow Q; Q \leftrightarrow R |] \implies P \leftrightarrow R$
 $\langle proof \rangle$

lemma *exII*:
 $P(a) \implies (!x. P(x) \implies x=a) \implies EX! x. P(x)$
 $\langle proof \rangle$

lemma *ex-ex1I*:

$EX\ x. P(x) \implies (!x\ y. [P(x); P(y)] \implies x=y) \implies EX!\ x. P(x)$
 $\langle proof \rangle$

lemma *ex1E*:

$EX!\ x. P(x) \implies (!x. [P(x); ALL\ y. P(y) \dashv\rightarrow y=x] \implies R) \implies R$
 $\langle proof \rangle$

$\langle ML \rangle$

lemma *conj-cong*:

assumes $P \dashv\rightarrow P'$
and $P' \implies Q \dashv\rightarrow Q'$
shows $(P \& Q) \dashv\rightarrow (P' \& Q')$
 $\langle proof \rangle$

lemma *conj-cong2*:

assumes $P \dashv\rightarrow P'$
and $P' \implies Q \dashv\rightarrow Q'$
shows $(Q \& P) \dashv\rightarrow (Q' \& P')$
 $\langle proof \rangle$

lemma *disj-cong*:

assumes $P \dashv\rightarrow P'$ **and** $Q \dashv\rightarrow Q'$
shows $(P | Q) \dashv\rightarrow (P' | Q')$
 $\langle proof \rangle$

lemma *imp-cong*:

assumes $P \dashv\rightarrow P'$
and $P' \implies Q \dashv\rightarrow Q'$
shows $(P \dashv\rightarrow Q) \dashv\rightarrow (P' \dashv\rightarrow Q')$
 $\langle proof \rangle$

lemma *iff-cong*: $[P \dashv\rightarrow P'; Q \dashv\rightarrow Q'] \implies (P \dashv\rightarrow Q) \dashv\rightarrow (P' \dashv\rightarrow Q')$
 $\langle proof \rangle$

lemma *not-cong*: $P \dashv\rightarrow P' \implies \sim P \dashv\rightarrow \sim P'$
 $\langle proof \rangle$

lemma *all-cong*:

assumes $!x. P(x) \dashv\rightarrow Q(x)$
shows $(ALL\ x. P(x)) \dashv\rightarrow (ALL\ x. Q(x))$

$\langle proof \rangle$

lemma *ex-cong*:

assumes $!!x. P(x) <-> Q(x)$

shows $(EX\ x. P(x)) <-> (EX\ x. Q(x))$

$\langle proof \rangle$

lemma *ex1-cong*:

assumes $!!x. P(x) <-> Q(x)$

shows $(EX!\ x. P(x)) <-> (EX!\ x. Q(x))$

$\langle proof \rangle$

lemma *sym*: $a=b ==> b=a$

$\langle proof \rangle$

lemma *trans*: $[| a=b; b=c |] ==> a=c$

$\langle proof \rangle$

lemma *not-sym*: $b \sim = a ==> a \sim = b$

$\langle proof \rangle$

lemma *def-imp-iff*: $(A == B) ==> A <-> B$

$\langle proof \rangle$

lemma *meta-eq-to-obj-eq*: $(A == B) ==> A = B$

$\langle proof \rangle$

lemma *meta-eq-to-iff*: $x==y ==> x<->y$

$\langle proof \rangle$

lemma *ssubst*: $[| b = a; P(a) |] ==> P(b)$

$\langle proof \rangle$

lemma *ex1-equalsE*:

$[| EX!\ x. P(x); P(a); P(b) |] ==> a=b$

$\langle proof \rangle$

lemma *subst-context*: $[| a=b |] ==> t(a)=t(b)$

$\langle proof \rangle$

lemma *subst-context2*: $[[a=b; c=d]] ==> t(a,c)=t(b,d)$
 $\langle proof \rangle$

lemma *subst-context3*: $[[a=b; c=d; e=f]] ==> t(a,c,e)=t(b,d,f)$
 $\langle proof \rangle$

lemma *box-equals*: $[[a=b; a=c; b=d]] ==> c=d$
 $\langle proof \rangle$

lemma *simp-equals*: $[[a=c; b=d; c=d]] ==> a=b$
 $\langle proof \rangle$

lemma *pred1-cong*: $a=a' ==> P(a) <-> P(a')$
 $\langle proof \rangle$

lemma *pred2-cong*: $[[a=a'; b=b']] ==> P(a,b) <-> P(a',b')$
 $\langle proof \rangle$

lemma *pred3-cong*: $[[a=a'; b=b'; c=c']] ==> P(a,b,c) <-> P(a',b',c')$
 $\langle proof \rangle$

$\langle ML \rangle$

lemma *eq-cong*: $[[a = a'; b = b']] ==> a = b <-> a' = b'$
 $\langle proof \rangle$

lemma *conj-impE*:
assumes *major*: $(P \& Q) \dashrightarrow S$
and *r*: $P \dashrightarrow (Q \dashrightarrow S) ==> R$
shows *R*
 $\langle proof \rangle$

lemma *disj-impE*:
assumes *major*: $(P | Q) \dashrightarrow S$
and *r*: $[[P \dashrightarrow S; Q \dashrightarrow S]] ==> R$
shows *R*
 $\langle proof \rangle$

lemma *imp-impE*:
assumes *major*: $(P \dashrightarrow Q) \dashrightarrow S$
and *r1*: $\llbracket P; Q \dashrightarrow S \rrbracket \implies Q$
and *r2*: $S \implies R$
shows *R*
 $\langle proof \rangle$

lemma *not-impE*:
 $\sim P \dashrightarrow S \implies (P \implies False) \implies (S \implies R) \implies R$
 $\langle proof \rangle$

lemma *iff-impE*:
assumes *major*: $(P <-> Q) \dashrightarrow S$
and *r1*: $\llbracket P; Q \dashrightarrow S \rrbracket \implies Q$
and *r2*: $\llbracket Q; P \dashrightarrow S \rrbracket \implies P$
and *r3*: $S \implies R$
shows *R*
 $\langle proof \rangle$

lemma *all-impE*:
assumes *major*: $(ALL\ x.\ P(x)) \dashrightarrow S$
and *r1*: $!!x.\ P(x)$
and *r2*: $S \implies R$
shows *R*
 $\langle proof \rangle$

lemma *ex-impE*:
assumes *major*: $(EX\ x.\ P(x)) \dashrightarrow S$
and *r*: $P(x) \dashrightarrow S \implies R$
shows *R*
 $\langle proof \rangle$

lemma *disj-imp-disj*:
 $P|Q \implies (P \implies R) \implies (Q \implies S) \implies R|S$
 $\langle proof \rangle$

$\langle ML \rangle$

lemma *thin-refl*: $!!X.\ \llbracket x=x; PROP\ W \rrbracket \implies PROP\ W\ \langle proof \rangle$

$\langle ML \rangle$

1.3 Intuitionistic Reasoning

lemma *impE'*:
 assumes 1: $P \dashv\vdash Q$
 and 2: $Q \implies R$
 and 3: $P \dashv\vdash Q \implies P$
 shows R
 $\langle proof \rangle$

lemma *allE'*:
 assumes 1: $\text{ALL } x. P(x)$
 and 2: $P(x) \implies \text{ALL } x. P(x) \implies Q$
 shows Q
 $\langle proof \rangle$

lemma *notE'*:
 assumes 1: $\sim P$
 and 2: $\sim P \implies P$
 shows R
 $\langle proof \rangle$

lemmas [*Pure.elim!*] = *disjE iffE FalseE conjE exE*
 and [*Pure.intro!*] = *iffI conjI impI TrueI notI allI refl*
 and [*Pure.elim 2*] = *allE notE' impE'*
 and [*Pure.intro*] = *exI disjI2 disjI1*

$\langle ML \rangle$

lemma *iff-not-sym*: $\sim (Q \dashv\vdash P) \implies \sim (P \dashv\vdash Q)$
 $\langle proof \rangle$

lemmas [*sym*] = *sym iff-sym not-sym iff-not-sym*
 and [*Pure.elim?*] = *iffD1 iffD2 impE*

lemma *eq-commute*: $a=b \dashv\vdash b=a$
 $\langle proof \rangle$

1.4 Atomizing meta-level rules

lemma *atomize-all* [*atomize*]: $(!!x. P(x)) \implies \text{Trueprop } (\text{ALL } x. P(x))$
 $\langle proof \rangle$

lemma *atomize-imp* [*atomize*]: $(A \implies B) \implies \text{Trueprop } (A \dashv\vdash B)$
 $\langle proof \rangle$

lemma *atomize-eq* [*atomize*]: $(x == y) \implies \text{Trueprop } (x = y)$
 $\langle proof \rangle$

lemma *atomize-iff* [*atomize*]: $(A == B) == \text{Trueprop } (A <-> B)$
 $\langle \text{proof} \rangle$

lemma *atomize-conj* [*atomize*]:
includes *meta-conjunction-syntax*
shows $(A \&\& B) == \text{Trueprop } (A \& B)$
 $\langle \text{proof} \rangle$

lemmas [*symmetric, rulify*] = *atomize-all atomize-imp*
and [*symmetric, defn*] = *atomize-all atomize-imp atomize-eq atomize-iff*

1.5 Atomizing elimination rules

$\langle ML \rangle$

lemma *atomize-exL*[*atomize-elim*]: $(!!x. P(x) ==> Q) == ((EX x. P(x)) ==> Q)$
 $\langle \text{proof} \rangle$

lemma *atomize-conjL*[*atomize-elim*]: $(A ==> B ==> C) == (A \& B ==> C)$
 $\langle \text{proof} \rangle$

lemma *atomize-disjL*[*atomize-elim*]: $((A ==> C) ==> (B ==> C) ==> C) == ((A | B ==> C) ==> C)$
 $\langle \text{proof} \rangle$

lemma *atomize-elimL*[*atomize-elim*]: $(!!B. (A ==> B) ==> B) == \text{Trueprop}(A)$
 $\langle \text{proof} \rangle$

1.6 Calculational rules

lemma *forw-subst*: $a = b ==> P(b) ==> P(a)$
 $\langle \text{proof} \rangle$

lemma *back-subst*: $P(a) ==> a = b ==> P(b)$
 $\langle \text{proof} \rangle$

Note that this list of rules is in reverse order of priorities.

lemmas *basic-trans-rules* [*trans*] =
forw-subst
back-subst
rev-mp
mp
trans

1.7 “Let” declarations

nonterminals *letbinds letbind*

constdefs

$Let :: ['a::\{\}, 'a ==> 'b] ==> ('b::\{\})$
 $Let(s, f) == f(s)$

syntax

$-bind :: [pttrn, 'a] ==> letbind \quad ((2- =/ -) 10)$
 $\quad \quad \quad :: letbind ==> letbinds \quad (-)$
 $-binds :: [letbind, letbinds] ==> letbinds \quad (-;/ -)$
 $-Let :: [letbinds, 'a] ==> 'a \quad ((let (-)/ in (-)) 10)$

translations

$-Let(-binds(b, bs), e) == -Let(b, -Let(bs, e))$
 $let\ x = a\ in\ e \quad ==\ Let(a, \%x. e)$

lemma LetI:

assumes $!!x. x=t ==> P(u(x))$
shows $P(let\ x=t\ in\ u(x))$
 $\langle proof \rangle$

1.8 Intuitionistic simplification rules

lemma conj-simps:

$P \ \& \ True <-> P$
 $True \ \& \ P <-> P$
 $P \ \& \ False <-> False$
 $False \ \& \ P <-> False$
 $P \ \& \ P <-> P$
 $P \ \& \ P \ \& \ Q <-> P \ \& \ Q$
 $P \ \& \ \sim P <-> False$
 $\sim P \ \& \ P <-> False$
 $(P \ \& \ Q) \ \& \ R <-> P \ \& \ (Q \ \& \ R)$
 $\langle proof \rangle$

lemma disj-simps:

$P \mid True <-> True$
 $True \mid P <-> True$
 $P \mid False <-> P$
 $False \mid P <-> P$
 $P \mid P <-> P$
 $P \mid P \mid Q <-> P \mid Q$
 $(P \mid Q) \mid R <-> P \mid (Q \mid R)$
 $\langle proof \rangle$

lemma not-simps:

$\sim(P \mid Q) <-> \sim P \ \& \ \sim Q$
 $\sim False <-> True$
 $\sim True <-> False$
 $\langle proof \rangle$

lemma *imp-simps*:

$(P \multimap \text{False}) \leftrightarrow \sim P$
 $(P \multimap \text{True}) \leftrightarrow \text{True}$
 $(\text{False} \multimap P) \leftrightarrow \text{True}$
 $(\text{True} \multimap P) \leftrightarrow P$
 $(P \multimap P) \leftrightarrow \text{True}$
 $(P \multimap \sim P) \leftrightarrow \sim P$
 $\langle \text{proof} \rangle$

lemma *iff-simps*:

$(\text{True} \leftrightarrow P) \leftrightarrow P$
 $(P \leftrightarrow \text{True}) \leftrightarrow P$
 $(P \leftrightarrow P) \leftrightarrow \text{True}$
 $(\text{False} \leftrightarrow P) \leftrightarrow \sim P$
 $(P \leftrightarrow \text{False}) \leftrightarrow \sim P$
 $\langle \text{proof} \rangle$

lemma *quant-simps*:

$!!P. (\text{ALL } x. P) \leftrightarrow P$
 $(\text{ALL } x. x=t \multimap P(x)) \leftrightarrow P(t)$
 $(\text{ALL } x. t=x \multimap P(x)) \leftrightarrow P(t)$
 $!!P. (\text{EX } x. P) \leftrightarrow P$
 $\text{EX } x. x=t$
 $\text{EX } x. t=x$
 $(\text{EX } x. x=t \ \& \ P(x)) \leftrightarrow P(t)$
 $(\text{EX } x. t=x \ \& \ P(x)) \leftrightarrow P(t)$
 $\langle \text{proof} \rangle$

lemma *distrib-simps*:

$P \ \& \ (Q \mid R) \leftrightarrow P \ \& \ Q \mid P \ \& \ R$
 $(Q \mid R) \ \& \ P \leftrightarrow Q \ \& \ P \mid R \ \& \ P$
 $(P \mid Q \multimap R) \leftrightarrow (P \multimap R) \ \& \ (Q \multimap R)$
 $\langle \text{proof} \rangle$

Conversion into rewrite rules

lemma *P-iff-F*: $\sim P \implies (P \leftrightarrow \text{False})$ $\langle \text{proof} \rangle$

lemma *iff-reflection-F*: $\sim P \implies (P == \text{False})$ $\langle \text{proof} \rangle$

lemma *P-iff-T*: $P \implies (P \leftrightarrow \text{True})$ $\langle \text{proof} \rangle$

lemma *iff-reflection-T*: $P \implies (P == \text{True})$ $\langle \text{proof} \rangle$

More rewrite rules

lemma *conj-commute*: $P \ \& \ Q \leftrightarrow Q \ \& \ P$ $\langle \text{proof} \rangle$

lemma *conj-left-commute*: $P \ \& \ (Q \ \& \ R) \leftrightarrow Q \ \& \ (P \ \& \ R)$ $\langle \text{proof} \rangle$

lemmas *conj-comms* = *conj-commute conj-left-commute*

lemma *disj-commute*: $P \mid Q \leftrightarrow Q \mid P$ $\langle \text{proof} \rangle$

lemma *disj-left-commute*: $P|(Q|R) <-> Q|(P|R)$ *<proof>*
lemmas *disj-comms* = *disj-commute disj-left-commute*

lemma *conj-disj-distribL*: $P\&(Q|R) <-> (P\&Q | P\&R)$ *<proof>*
lemma *conj-disj-distribR*: $(P|Q)\&R <-> (P\&R | Q\&R)$ *<proof>*

lemma *disj-conj-distribL*: $P|(Q\&R) <-> (P|Q) \& (P|R)$ *<proof>*
lemma *disj-conj-distribR*: $(P\&Q)|R <-> (P|R) \& (Q|R)$ *<proof>*

lemma *imp-conj-distrib*: $(P \dashrightarrow (Q\&R)) <-> (P \dashrightarrow Q) \& (P \dashrightarrow R)$ *<proof>*
lemma *imp-conj*: $((P\&Q) \dashrightarrow R) <-> (P \dashrightarrow (Q \dashrightarrow R))$ *<proof>*
lemma *imp-disj*: $(P|Q \dashrightarrow R) <-> (P \dashrightarrow R) \& (Q \dashrightarrow R)$ *<proof>*

lemma *de-Morgan-disj*: $(\sim(P | Q)) <-> (\sim P \& \sim Q)$ *<proof>*

lemma *not-ex*: $(\sim (EX\ x.\ P(x))) <-> (ALL\ x.\ \sim P(x))$ *<proof>*
lemma *imp-ex*: $((EX\ x.\ P(x)) \dashrightarrow Q) <-> (ALL\ x.\ P(x) \dashrightarrow Q)$ *<proof>*

lemma *ex-disj-distrib*:
 $(EX\ x.\ P(x) | Q(x)) <-> ((EX\ x.\ P(x)) | (EX\ x.\ Q(x)))$ *<proof>*

lemma *all-conj-distrib*:
 $(ALL\ x.\ P(x) \& Q(x)) <-> ((ALL\ x.\ P(x)) \& (ALL\ x.\ Q(x)))$ *<proof>*

1.9 Legacy ML bindings

<ML>

end

2 Classical first-order logic

theory *FOL*
imports *IFOL*
uses
 $\sim\sim$ */src/Provers/classical.ML*
 $\sim\sim$ */src/Provers/blast.ML*
 $\sim\sim$ */src/Provers/clasimp.ML*
 $\sim\sim$ */src/Tools/induct.ML*
(cladata.ML)
(blastdata.ML)
(simpdata.ML)
begin

2.1 The classical axiom

axioms
classical: $(\sim P ==> P) ==> P$

2.2 Lemmas and proof tools

lemma *ccontr*: $(\neg P \implies \text{False}) \implies P$
<proof>

lemma *disjCI*: $(\sim Q \implies P) \implies P \mid Q$
<proof>

lemma *ex-classical*:
 assumes *r*: $\sim (EX\ x. P(x)) \implies P(a)$
 shows $EX\ x. P(x)$
<proof>

lemma *exCI*:
 assumes *r*: $ALL\ x. \sim P(x) \implies P(a)$
 shows $EX\ x. P(x)$
<proof>

lemma *excluded-middle*: $\sim P \mid P$
<proof>

<ML>

lemma *case-split-thm*:
 assumes *r1*: $P \implies Q$
 and *r2*: $\sim P \implies Q$
 shows Q
<proof>

lemmas *case-split* = *case-split-thm* [*case-names True False*]

<ML>

lemma *impCE*:
 assumes *major*: $P \dashv\dashv Q$
 and *r1*: $\sim P \implies R$
 and *r2*: $Q \implies R$
 shows R
<proof>

lemma *impCE'*:
assumes *major*: $P \dashrightarrow Q$
and *r1*: $Q \implies R$
and *r2*: $\sim P \implies R$
shows R
 $\langle proof \rangle$

lemma *notnotD*: $\sim\sim P \implies P$
 $\langle proof \rangle$

lemma *contrapos2*: $[\![\ Q; \sim P \implies \sim Q\]\!] \implies P$
 $\langle proof \rangle$

lemma *iffCE*:
assumes *major*: $P <-> Q$
and *r1*: $[\![\ P; Q\]\!] \implies R$
and *r2*: $[\![\ \sim P; \sim Q\]\!] \implies R$
shows R
 $\langle proof \rangle$

lemma *alt-ex1E*:
assumes *major*: $EX! x. P(x)$
and *r*: $!!x. [\![\ P(x); ALL\ y\ y'. P(y) \ \&\ P(y') \dashrightarrow y=y'\]\!] \implies R$
shows R
 $\langle proof \rangle$

lemma *imp-elim*: $P \dashrightarrow Q \implies (\sim R \implies P) \implies (Q \implies R) \implies R$
 $\langle proof \rangle$

lemma *swap*: $\sim P \implies (\sim R \implies P) \implies R$
 $\langle proof \rangle$

$\langle ML \rangle$

lemma *ex1-functional*: $[\![\ EX! z. P(a,z); P(a,b); P(a,c)\]\!] \implies b = c$
 $\langle proof \rangle$

lemma *True-implies-equals*: $(True \implies PROP\ P) == PROP\ P$
 $\langle proof \rangle$

lemma *uncurry*: $P \multimap Q \multimap R \implies P \& Q \multimap R$
 $\langle proof \rangle$

lemma *iff-allI*: $(!!x. P(x) \iff Q(x)) \implies (ALL\ x. P(x) \iff ALL\ x. Q(x))$
 $\langle proof \rangle$

lemma *iff-exI*: $(!!x. P(x) \iff Q(x)) \implies (EX\ x. P(x) \iff EX\ x. Q(x))$
 $\langle proof \rangle$

lemma *all-comm*: $(ALL\ x\ y. P(x,y)) \iff (ALL\ y\ x. P(x,y))$ $\langle proof \rangle$

lemma *ex-comm*: $(EX\ x\ y. P(x,y)) \iff (EX\ y\ x. P(x,y))$ $\langle proof \rangle$

lemma *cases-simp*: $(P \multimap Q) \& (\sim P \multimap Q) \iff Q$ $\langle proof \rangle$

lemma *int-ex-simps*:
 $!!P\ Q. (EX\ x. P(x) \& Q) \iff (EX\ x. P(x)) \& Q$
 $!!P\ Q. (EX\ x. P \& Q(x)) \iff P \& (EX\ x. Q(x))$
 $!!P\ Q. (EX\ x. P(x) \mid Q) \iff (EX\ x. P(x)) \mid Q$
 $!!P\ Q. (EX\ x. P \mid Q(x)) \iff P \mid (EX\ x. Q(x))$
 $\langle proof \rangle$

lemma *cla-ex-simps*:
 $!!P\ Q. (EX\ x. P(x) \multimap Q) \iff (ALL\ x. P(x)) \multimap Q$
 $!!P\ Q. (EX\ x. P \multimap Q(x)) \iff P \multimap (EX\ x. Q(x))$
 $\langle proof \rangle$

lemmas *ex-simps = int-ex-simps cla-ex-simps*

lemma *int-all-simps*:
 $!!P\ Q. (ALL\ x. P(x) \& Q) \iff (ALL\ x. P(x)) \& Q$
 $!!P\ Q. (ALL\ x. P \& Q(x)) \iff P \& (ALL\ x. Q(x))$
 $!!P\ Q. (ALL\ x. P(x) \multimap Q) \iff (EX\ x. P(x)) \multimap Q$
 $!!P\ Q. (ALL\ x. P \multimap Q(x)) \iff P \multimap (ALL\ x. Q(x))$
 $\langle proof \rangle$

lemma *cla-all-simps*:

!!P Q. (ALL x. P(x) | Q) <-> (ALL x. P(x)) | Q
 !!P Q. (ALL x. P | Q(x)) <-> P | (ALL x. Q(x))
 <proof>

lemmas *all-simps = int-all-simps cla-all-simps*

lemma *imp-disj1*: (P-->Q) | R <-> (P-->Q | R) <proof>

lemma *imp-disj2*: Q | (P-->R) <-> (P-->Q | R) <proof>

lemma *de-Morgan-conj*: (~ (P & Q)) <-> (~ P | ~ Q) <proof>

lemma *not-imp*: ~ (P --> Q) <-> (P & ~ Q) <proof>

lemma *not-iff*: ~ (P <-> Q) <-> (P <-> ~ Q) <proof>

lemma *not-all*: (~ (ALL x. P(x))) <-> (EX x. ~ P(x)) <proof>

lemma *imp-all*: ((ALL x. P(x)) --> Q) <-> (EX x. P(x) --> Q) <proof>

lemmas *meta-simps =*

triv-forall-equality

True-implies-equals

lemmas *IFOL-simps =*

refl [THEN P-iff-T] conj-simps disj-simps not-simps

imp-simps iff-simps quant-simps

lemma *notFalseI*: ~ False <proof>

lemma *cla-simps-misc*:

~ (P & Q) <-> ~ P | ~ Q

P | ~ P

~ P | P

~ ~ P <-> P

(~ P --> P) <-> P

(~ P <-> ~ Q) <-> (P <-> Q) <proof>

lemmas *cla-simps =*

de-Morgan-conj de-Morgan-disj imp-disj1 imp-disj2

not-imp not-all not-ex cases-simp cla-simps-misc

<ML>

2.3 Other simple lemmas

lemma *[simp]*: $((P \multimap R) \leftrightarrow (Q \multimap R)) \leftrightarrow ((P \leftrightarrow Q) \mid R)$
 $\langle proof \rangle$

lemma *[simp]*: $((P \multimap Q) \leftrightarrow (P \multimap R)) \leftrightarrow (P \multimap (Q \leftrightarrow R))$
 $\langle proof \rangle$

lemma *not-disj-iff-imp*: $\sim P \mid Q \leftrightarrow (P \multimap Q)$
 $\langle proof \rangle$

lemma *conj-mono*: $[[P1 \multimap Q1; P2 \multimap Q2]] \implies (P1 \& P2) \multimap (Q1 \& Q2)$
 $\langle proof \rangle$

lemma *disj-mono*: $[[P1 \multimap Q1; P2 \multimap Q2]] \implies (P1 \mid P2) \multimap (Q1 \mid Q2)$
 $\langle proof \rangle$

lemma *imp-mono*: $[[Q1 \multimap P1; P2 \multimap Q2]] \implies (P1 \multimap P2) \multimap (Q1 \multimap Q2)$
 $\langle proof \rangle$

lemma *imp-refl*: $P \multimap P$
 $\langle proof \rangle$

lemma *ex-mono*: $(!!x. P(x) \multimap Q(x)) \implies (EX x. P(x)) \multimap (EX x. Q(x))$
 $\langle proof \rangle$

lemma *all-mono*: $(!!x. P(x) \multimap Q(x)) \implies (ALL x. P(x)) \multimap (ALL x. Q(x))$
 $\langle proof \rangle$

2.4 Proof by cases and induction

Proper handling of non-atomic rule statements.

constdefs

induct-forall **where** *induct-forall*(P) == $\forall x. P(x)$
induct-implies **where** *induct-implies*(A, B) == $A \longrightarrow B$
induct-equal **where** *induct-equal*(x, y) == $x = y$
induct-conj **where** *induct-conj*(A, B) == $A \wedge B$

lemma *induct-forall-eq*: $(!!x. P(x)) == \text{Trueprop}(\text{induct-forall}(\lambda x. P(x)))$
 $\langle proof \rangle$

lemma *induct-implies-eq*: $(A \implies B) == \text{Trueprop}(\text{induct-implies}(A, B))$
 $\langle proof \rangle$

lemma *induct-equal-eq*: $(x == y) == \text{Trueprop}(\text{induct-equal}(x, y))$

```

    <proof>

lemma induct-conj-eq:
  includes meta-conjunction-syntax
  shows  $(A \ \&\& \ B) == \text{Trueprop}(\text{induct-conj}(A, B))$ 
  <proof>

lemmas induct-atomize = induct-forall-eq induct-implies-eq induct-equal-eq induct-conj-eq
lemmas induct-rulify [symmetric, standard] = induct-atomize
lemmas induct-rulify-fallback =
  induct-forall-def induct-implies-def induct-equal-def induct-conj-def

hide const induct-forall induct-implies induct-equal induct-conj

Method setup.

<ML>
declare case-split [cases type: o]

end

```