

# Machine Words in Isabelle/HOL

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## Abstract

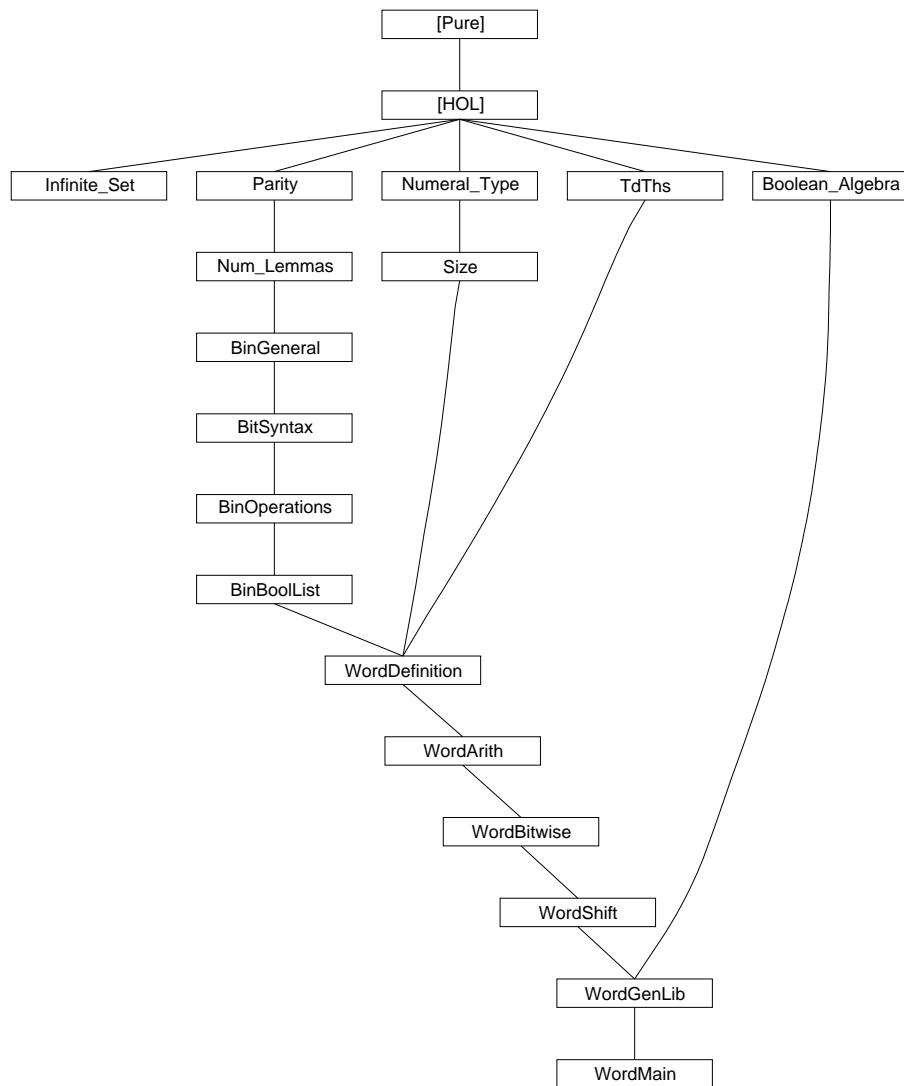
A formalisation of generic, fixed size machine words in Isabelle/HOL.  
An earlier version of this formalisation is described in [1].

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# 1 Numeral-Type: Numeral Syntax for Types

```
theory Numeral-Type
  imports ATP-Linkup
begin
```

## 1.1 Preliminary lemmas

```
lemma inj-Inl [simp]: inj-on Inl A
  <proof>
```

```
lemma inj-Inr [simp]: inj-on Inr A
  <proof>
```

```
lemma inj-Some [simp]: inj-on Some A
  <proof>
```

```
lemma card-Plus:
  [| finite A; finite B |] ==> card (A <+> B) = card A + card B
  <proof>
```

```
lemma (in type-definition) univ:
  UNIV = Abs ‘ A
  <proof>
```

```
lemma (in type-definition) card: card (UNIV :: 'b set) = card A
  <proof>
```

## 1.2 Cardinalities of types

```
syntax -type-card :: type => nat ((1CARD/(1'(-))))
```

```
translations CARD(t) => CONST card (UNIV :: t set)
```

```
<ML>
```

```
lemma card-unit: CARD(unit) = 1
  <proof>
```

```
lemma card-bool: CARD(bool) = 2
  <proof>
```

```
lemma card-prod: CARD('a::finite × 'b::finite) = CARD('a) * CARD('b)
  <proof>
```

```
lemma card-sum: CARD('a::finite + 'b::finite) = CARD('a) + CARD('b)
  <proof>
```

```
lemma card-option: CARD('a::finite option) = Suc CARD('a)
  <proof>
```

**lemma** *card-set*:  $CARD('a::finite\ set) = 2 \wedge CARD('a)$   
 $\langle proof \rangle$

### 1.3 Numeral Types

**typedef** (**open**) *num0* = *UNIV* :: *nat set*  $\langle proof \rangle$   
**typedef** (**open**) *num1* = *UNIV* :: *unit set*  $\langle proof \rangle$   
**typedef** (**open**) *'a bit0* = *UNIV* :: (*bool* \* *'a*) *set*  $\langle proof \rangle$   
**typedef** (**open**) *'a bit1* = *UNIV* :: (*bool* \* *'a*) *option set*  $\langle proof \rangle$

**instance** *num1* :: *finite*  
 $\langle proof \rangle$

**instance** *bit0* :: (*finite*) *finite*  
 $\langle proof \rangle$

**instance** *bit1* :: (*finite*) *finite*  
 $\langle proof \rangle$

**lemma** *card-num1*:  $CARD(num1) = 1$   
 $\langle proof \rangle$

**lemma** *card-bit0*:  $CARD('a::finite\ bit0) = 2 * CARD('a)$   
 $\langle proof \rangle$

**lemma** *card-bit1*:  $CARD('a::finite\ bit1) = Suc\ (2 * CARD('a))$   
 $\langle proof \rangle$

**lemma** *card-num0*:  $CARD\ (num0) = 0$   
 $\langle proof \rangle$

**lemmas** *card-univ-simps* [*simp*] =  
*card-unit*  
*card-bool*  
*card-prod*  
*card-sum*  
*card-option*  
*card-set*  
*card-num1*  
*card-bit0*  
*card-bit1*  
*card-num0*

### 1.4 Syntax

**syntax**  
*-NumeralType* :: *num-const* => *type* (-)  
*-NumeralType0* :: *type* (0)  
*-NumeralType1* :: *type* (1)

**translations**

-*NumeralType1* == (*type*) *num1*  
 -*NumeralType0* == (*type*) *num0*

⟨*ML*⟩

**1.5 Classes with at least 1 and 2**

Class *finite* already captures “at least 1”

**lemma** *zero-less-card-finite* [*simp*]:  
 $0 < \text{CARD}('a::\text{finite})$   
 ⟨*proof*⟩

**lemma** *one-le-card-finite* [*simp*]:  
 $\text{Suc } 0 \leq \text{CARD}('a::\text{finite})$   
 ⟨*proof*⟩

Class for cardinality “at least 2”

**class** *card2* = *finite* +  
**assumes** *two-le-card*:  $2 \leq \text{CARD}('a)$

**lemma** *one-less-card*:  $\text{Suc } 0 < \text{CARD}('a::\text{card2})$   
 ⟨*proof*⟩

**instance** *bit0* :: (*finite*) *card2*  
 ⟨*proof*⟩

**instance** *bit1* :: (*finite*) *card2*  
 ⟨*proof*⟩

**1.6 Examples**

**lemma**  $\text{CARD}(0) = 0$  ⟨*proof*⟩  
**lemma**  $\text{CARD}(17) = 17$  ⟨*proof*⟩

**end**

**2 Size: The len classes**

**theory** *Size*  
**imports** *Numeral-Type*  
**begin**

The aim of this is to allow any type as index type, but to provide a default instantiation for numeral types. This independence requires some duplication with the definitions in *Numeral-Type*.

```

class len0 = type +
  fixes len-of :: 'a itself ⇒ nat

```

Some theorems are only true on words with length greater 0.

```

class len = len0 +
  assumes len-gt-0 [iff]:  $0 < \text{len-of } \text{TYPE } ('a)$ 

```

```

instantiation num0 and num1 :: len0
begin

```

```

definition
  len-num0: len-of (x::num0 itself) = 0

```

```

definition
  len-num1: len-of (x::num1 itself) = 1

```

```

instance ⟨proof⟩

```

```

end

```

```

instantiation bit0 and bit1 :: (len0) len0
begin

```

```

definition
  len-bit0: len-of (x::'a::len0 bit0 itself) = 2 * len-of TYPE ('a)

```

```

definition
  len-bit1: len-of (x::'a::len0 bit1 itself) = 2 * len-of TYPE ('a) + 1

```

```

instance ⟨proof⟩

```

```

end

```

```

lemmas len-of-numeral-defs [simp] = len-num0 len-num1 len-bit0 len-bit1

```

```

instance num1 :: len ⟨proof⟩
instance bit0 :: (len) len ⟨proof⟩
instance bit1 :: (len0) len ⟨proof⟩
lemma len-of TYPE(17) = 17 ⟨proof⟩
lemma len-of TYPE(0) = 0 ⟨proof⟩
lemma len-of TYPE('a::len0) = x
  ⟨proof⟩

```

```

end

```



### 3 Num-Lemmas: Useful Numerical Lemmas

```
theory Num-Lemmas
imports Main Parity
begin
```

```
lemma contentsI:  $y = \{x\} \implies \text{contents } y = x$ 
  <proof>
```

```
lemmas split-split = prod.split [unfolded prod-case-split]
lemmas split-split-asm = prod.split-asm [unfolded prod-case-split]
lemmas split.splits = split-split split-split-asm
```

```
lemmas funpow-0 = funpow.simps(1)
lemmas funpow-Suc = funpow.simps(2)
```

```
lemma nonemptyE:  $S \sim \{ \} \implies (!x. x : S \implies R) \implies R$ 
  <proof>
```

```
lemma gt-or-eq-0:  $0 < y \vee 0 = (y :: \text{nat})$  <proof>
```

```
declare iszero-0 [iff]
```

```
lemmas xtr1 = xtrans(1)
lemmas xtr2 = xtrans(2)
lemmas xtr3 = xtrans(3)
lemmas xtr4 = xtrans(4)
lemmas xtr5 = xtrans(5)
lemmas xtr6 = xtrans(6)
lemmas xtr7 = xtrans(7)
lemmas xtr8 = xtrans(8)
```

```
lemmas nat-simps = diff-add-inverse2 diff-add-inverse
lemmas nat-iffs = le-add1 le-add2
```

```
lemma sum-imp-diff:  $j = k + i \implies j - i = (k :: \text{nat})$ 
  <proof>
```

```
lemma nobm1:
   $0 < (\text{number-of } w :: \text{nat}) \implies$ 
   $\text{number-of } w - (1 :: \text{nat}) = \text{number-of } (\text{Int.pred } w)$ 
  <proof>
```

```
lemma of-int-power:
   $\text{of-int } (a \wedge n) = (\text{of-int } a \wedge n :: 'a :: \{\text{recpower, comm-ring-1}\})$ 
  <proof>
```

```
lemma zless2:  $0 < (2 :: \text{int})$ 
  <proof>
```

**lemmas**  $zless2p$  [simp] =  $zless2$  [THEN zero-less-power]

**lemmas**  $zle2p$  [simp] =  $zless2p$  [THEN order-less-imp-le]

**lemmas**  $pos-mod-sign2$  =  $zless2$  [THEN pos-mod-sign [where  $b = 2::int$ ]]

**lemmas**  $pos-mod-bound2$  =  $zless2$  [THEN pos-mod-bound [where  $b = 2::int$ ]]

— the inverse(s) of *number-of*

**lemma**  $nmod2$ :  $n \bmod (2::int) = 0 \mid n \bmod 2 = 1$

*<proof>*

**lemma**  $emep1$ :

*even*  $n ==> \text{even } d ==> 0 \leq d ==> (n + 1) \bmod (d :: int) = (n \bmod d) + 1$

*<proof>*

**lemmas**  $eme1p$  =  $emep1$  [simplified add-commute]

**lemma**  $le\text{-}diff\text{-}eq'$ :  $(a \leq c - b) = (b + a \leq (c::int))$

*<proof>*

**lemma**  $less\text{-}diff\text{-}eq'$ :  $(a < c - b) = (b + a < (c::int))$

*<proof>*

**lemma**  $diff\text{-}le\text{-}eq'$ :  $(a - b \leq c) = (a \leq b + (c::int))$

*<proof>*

**lemma**  $diff\text{-}less\text{-}eq'$ :  $(a - b < c) = (a < b + (c::int))$

*<proof>*

**lemmas**  $m1mod2k$  =  $zless2p$  [THEN  $zmod\text{-}minus1$ ]

**lemmas**  $m1mod22k$  =  $mult\text{-}pos\text{-}pos$  [OF  $zless2$   $zless2p$ , THEN  $zmod\text{-}minus1$ ]

**lemmas**  $p1mod22k'$  =  $zless2p$  [THEN order-less-imp-le, THEN pos-zmod-mult-2]

**lemmas**  $z1pmod2'$  = zero-le-one [THEN pos-zmod-mult-2, simplified]

**lemmas**  $z1pdiv2'$  = zero-le-one [THEN pos-zdiv-mult-2, simplified]

**lemma**  $p1mod22k$ :

$(2 * b + 1) \bmod (2 * 2^n) = 2 * (b \bmod 2^n) + (1::int)$

*<proof>*

**lemma**  $z1pmod2$ :

$(2 * b + 1) \bmod 2 = (1::int)$

*<proof>*

**lemma**  $z1pdiv2$ :

$(2 * b + 1) \bmod 2 = (b::int)$

*<proof>*

**lemmas** *zdiv-le-dividend* = *xtr3* [*OF* *zdiv-1* [*symmetric*] *zdiv-mono2*,  
*simplified int-one-le-iff-zero-less, simplified, standard*]

**lemma** *axxbyy*:

$a + m + m = b + n + n \implies (a = 0 \mid a = 1) \implies (b = 0 \mid b = 1) \implies$   
 $a = b \ \& \ m = (n :: \text{int})$   
 $\langle \text{proof} \rangle$

**lemma** *axxmod2*:

$(1 + x + x) \bmod 2 = (1 :: \text{int}) \ \& \ (0 + x + x) \bmod 2 = (0 :: \text{int})$   
 $\langle \text{proof} \rangle$

**lemma** *axxdiv2*:

$(1 + x + x) \text{ div } 2 = (x :: \text{int}) \ \& \ (0 + x + x) \text{ div } 2 = (x :: \text{int})$   
 $\langle \text{proof} \rangle$

**lemmas** *iszero-minus* = *trans* [*THEN* *trans*,

*OF iszero-def neg-equal-0-iff-equal iszero-def* [*symmetric*], *standard*]

**lemmas** *zadd-diff-inverse* = *trans* [*OF* *diff-add-cancel* [*symmetric*] *add-commute*,  
*standard*]

**lemmas** *add-diff-cancel2* = *add-commute* [*THEN* *diff-eq-eq* [*THEN* *iffD2*], *standard*]

**lemma** *zmod-uminus*:  $-(a :: \text{int}) \bmod b \bmod b = -a \bmod b$   
 $\langle \text{proof} \rangle$

**lemma** *zmod-zsub-distrib*:  $((a :: \text{int}) - b) \bmod c = (a \bmod c - b \bmod c) \bmod c$   
 $\langle \text{proof} \rangle$

**lemma** *zmod-zsub-right-eq*:  $((a :: \text{int}) - b) \bmod c = (a - b \bmod c) \bmod c$   
 $\langle \text{proof} \rangle$

**lemma** *zmod-zsub-left-eq*:  $((a :: \text{int}) - b) \bmod c = (a \bmod c - b) \bmod c$   
 $\langle \text{proof} \rangle$

**lemma** *zmod-zsub-self* [*simp*]:

$((b :: \text{int}) - a) \bmod a = b \bmod a$   
 $\langle \text{proof} \rangle$

**lemma** *zmod-zmult1-eq-rev*:

$b * a \bmod c = b \bmod c * a \bmod (c :: \text{int})$   
 $\langle \text{proof} \rangle$

**lemmas** *rdmods* [*symmetric*] = *zmod-uminus* [*symmetric*]

*zmod-zsub-left-eq* *zmod-zsub-right-eq* *zmod-zadd-left-eq*

*zmod-zadd-right-eq* *zmod-zmult1-eq* *zmod-zmult1-eq-rev*

**lemma** *mod-plus-right*:

$((a + x) \bmod m = (b + x) \bmod m) = (a \bmod m = b \bmod m :: \text{nat})$   
 $\langle \text{proof} \rangle$

**lemma** *nat-minus-mod*:  $(n - n \bmod m) \bmod m = (0 :: \text{nat})$

$\langle \text{proof} \rangle$

**lemmas** *nat-minus-mod-plus-right* = *trans* [*OF* *nat-minus-mod mod-0* [*symmetric*],  
*THEN mod-plus-right* [*THEN iffD2*], *standard*, *simplified*]

**lemmas** *push-mods'* = *zmod-zadd1-eq* [*standard*]  
*zmod-zmult-distrib* [*standard*] *zmod-zsub-distrib* [*standard*]  
*zmod-uminus* [*symmetric*, *standard*]

**lemmas** *push-mods* = *push-mods'* [*THEN eq-reflection*, *standard*]

**lemmas** *pull-mods* = *push-mods* [*symmetric*] *rdmods* [*THEN eq-reflection*, *standard*]

**lemmas** *mod-simps* =  
*zmod-zmult-self1* [*THEN eq-reflection*] *zmod-zmult-self2* [*THEN eq-reflection*]  
*mod-mod-trivial* [*THEN eq-reflection*]

**lemma** *nat-mod-eq*:

$!!b. b < n ==> a \bmod n = b \bmod n ==> a \bmod n = (b :: \text{nat})$   
 $\langle \text{proof} \rangle$

**lemmas** *nat-mod-eq'* = *refl* [*THEN* [2] *nat-mod-eq*]

**lemma** *nat-mod-lem*:

$(0 :: \text{nat}) < n ==> b < n = (b \bmod n = b)$   
 $\langle \text{proof} \rangle$

**lemma** *mod-nat-add*:

$(x :: \text{nat}) < z ==> y < z ==>$   
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$   
 $\langle \text{proof} \rangle$

**lemma** *mod-nat-sub*:

$(x :: \text{nat}) < z ==> (x - y) \bmod z = x - y$   
 $\langle \text{proof} \rangle$

**lemma** *int-mod-lem*:

$(0 :: \text{int}) < n ==> (0 \leq b \ \& \ b < n) = (b \bmod n = b)$   
 $\langle \text{proof} \rangle$

**lemma** *int-mod-eq*:

$(0 :: \text{int}) \leq b ==> b < n ==> a \bmod n = b \bmod n ==> a \bmod n = b$   
 $\langle \text{proof} \rangle$

**lemmas** *int-mod-eq'* = *refl* [*THEN* [3] *int-mod-eq*]

**lemma** *int-mod-le*:  $0 \leq a \implies 0 < (n :: \text{int}) \implies a \bmod n \leq a$   
 ⟨proof⟩

**lemma** *int-mod-le'*:  $0 \leq b - n \implies 0 < (n :: \text{int}) \implies b \bmod n \leq b - n$   
 ⟨proof⟩

**lemma** *int-mod-ge*:  $a < n \implies 0 < (n :: \text{int}) \implies a \leq a \bmod n$   
 ⟨proof⟩

**lemma** *int-mod-ge'*:  $b < 0 \implies 0 < (n :: \text{int}) \implies b + n \leq b \bmod n$   
 ⟨proof⟩

**lemma** *mod-add-if-z*:  
 $(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$   
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$   
 ⟨proof⟩

**lemma** *mod-sub-if-z*:  
 $(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$   
 $(x - y) \bmod z = (\text{if } y \leq x \text{ then } x - y \text{ else } x - y + z)$   
 ⟨proof⟩

**lemmas** *zmde* = *zmod-zdiv-equality* [THEN *diff-eq-eq* [THEN *iffD2*], *symmetric*]  
**lemmas** *mcl* = *mult-cancel-left* [THEN *iffD1*, THEN *make-pos-rule*]

**lemma** *zdiv-mult-self*:  $m \sim = (0 :: \text{int}) \implies (a + m * n) \text{ div } m = a \text{ div } m + n$   
 ⟨proof⟩

**lemma** *eqne*:  $\text{equiv } A \ r \implies X : A // r \implies X \sim = \{\}$   
 ⟨proof⟩

**lemmas** *Rep-Integ-ne* = *Integ.Rep-Integ*  
 [THEN *equiv-intrel* [THEN *eqne*, *simplified Integ-def* [symmetric]], *standard*]

**lemmas** *riq* = *Integ.Rep-Integ* [*simplified Integ-def*]  
**lemmas** *intrel-refl* = *refl* [THEN *equiv-intrel-iff* [THEN *iffD1*], *standard*]  
**lemmas** *Rep-Integ-equiv* = *quotient-eq-iff*  
 [OF *equiv-intrel riq riq*, *simplified Integ.Rep-Integ-inject*, *standard*]  
**lemmas** *Rep-Integ-same* =  
*Rep-Integ-equiv* [THEN *intrel-refl* [THEN *rev-iffD2*], *standard*]

**lemma** *RI-int*:  $(a, 0) : \text{Rep-Integ } (\text{int } a)$   
 ⟨proof⟩

**lemmas** *RI-intrel* [*simp*] = *UNIV-I* [THEN *quotientI*,  
 THEN *Integ.Abs-Integ-inverse* [*simplified Integ-def*], *standard*]

**lemma** *RI-minus*:  $(a, b) : \text{Rep-Integ } x \implies (b, a) : \text{Rep-Integ } (-x)$   
 $\langle \text{proof} \rangle$

**lemma** *RI-add*:  
 $(a, b) : \text{Rep-Integ } x \implies (c, d) : \text{Rep-Integ } y \implies$   
 $(a + c, b + d) : \text{Rep-Integ } (x + y)$   
 $\langle \text{proof} \rangle$

**lemma** *mem-same*:  $a : S \implies a = b \implies b : S$   
 $\langle \text{proof} \rangle$

**lemma** *RI-eq-diff'*:  $(a, b) : \text{Rep-Integ } (\text{int } a - \text{int } b)$   
 $\langle \text{proof} \rangle$

**lemma** *RI-eq-diff*:  $((a, b) : \text{Rep-Integ } x) = (\text{int } a - \text{int } b = x)$   
 $\langle \text{proof} \rangle$

**lemma** *mod-power-lem*:  
 $a > 1 \implies a^n \bmod a^m = (\text{if } m \leq n \text{ then } 0 \text{ else } (a :: \text{int})^n)$   
 $\langle \text{proof} \rangle$

**lemma** *min-pm* [simp]:  $\min a b + (a - b) = (a :: \text{nat})$   
 $\langle \text{proof} \rangle$

**lemmas** *min-pm1* [simp] = trans [OF add-commute min-pm]

**lemma** *rev-min-pm* [simp]:  $\min b a + (a - b) = (a :: \text{nat})$   
 $\langle \text{proof} \rangle$

**lemmas** *rev-min-pm1* [simp] = trans [OF add-commute rev-min-pm]

**lemma** *pl-pl-rels*:  
 $a + b = c + d \implies$   
 $a \geq c \ \& \ b \leq d \mid a \leq c \ \& \ b \geq d \implies (d :: \text{nat})$   
 $\langle \text{proof} \rangle$

**lemmas** *pl-pl-rels'* = add-commute [THEN [2] trans, THEN pl-pl-rels]

**lemma** *minus-eq*:  $(m - k = m) = (k = 0 \mid m = (0 :: \text{nat}))$   
 $\langle \text{proof} \rangle$

**lemma** *pl-pl-mm*:  $(a :: \text{nat}) + b = c + d \implies a - c = d - b$   
 $\langle \text{proof} \rangle$

**lemmas** *pl-pl-mm'* = add-commute [THEN [2] trans, THEN pl-pl-mm]

**lemma** *min-minus* [simp]:  $\min m (m - k) = (m - k :: \text{nat})$

$\langle \text{proof} \rangle$

**lemmas** *min-minus'* [*simp*] = *trans* [*OF min-max.inf-commute min-minus*]

**lemma** *nat-no-eq-iff*:

$(\text{number-of } b :: \text{int}) \geq 0 \implies (\text{number-of } c :: \text{int}) \geq 0 \implies$   
 $(\text{number-of } b = (\text{number-of } c :: \text{nat})) = (b = c)$   
 $\langle \text{proof} \rangle$

**lemmas** *dme* = *box-equals* [*OF div-mod-equality add-0-right add-0-right*]

**lemmas** *dtle* = *xtr3* [*OF dme* [*symmetric*] *le-add1*]

**lemmas** *th2* = *order-trans* [*OF order-refl* [*THEN* [*2*] *mult-le-mono*] *dtle*]

**lemma** *td-gal*:

$0 < c \implies (a \geq b * c) = (a \text{ div } c \geq (b :: \text{nat}))$   
 $\langle \text{proof} \rangle$

**lemmas** *td-gal-lt* = *td-gal* [*simplified not-less* [*symmetric*], *simplified*]

**lemma** *div-mult-le*:  $(a :: \text{nat}) \text{ div } b * b \leq a$

$\langle \text{proof} \rangle$

**lemmas** *sdl* = *split-div-lemma* [*THEN iffD1*, *symmetric*]

**lemma** *given-quot*:  $f > (0 :: \text{nat}) \implies (f * l + (f - 1)) \text{ div } f = l$

$\langle \text{proof} \rangle$

**lemma** *given-quot-alt*:  $f > (0 :: \text{nat}) \implies (l * f + f - \text{Suc } 0) \text{ div } f = l$

$\langle \text{proof} \rangle$

**lemma** *diff-mod-le*:  $(a :: \text{nat}) < d \implies b \text{ dvd } d \implies a - a \text{ mod } b \leq d - b$

$\langle \text{proof} \rangle$

**lemma** *less-le-mult'*:

$w * c < b * c \implies 0 \leq c \implies (w + 1) * c \leq b * (c :: \text{int})$   
 $\langle \text{proof} \rangle$

**lemmas** *less-le-mult* = *less-le-mult'* [*simplified left-distrib*, *simplified*]

**lemmas** *less-le-mult-minus* = *iffD2* [*OF le-diff-eq less-le-mult*,  
*simplified left-diff-distrib*, *standard*]

**lemma** *lrlem'*:

**assumes** *d*:  $(i :: \text{nat}) \leq j \vee m < j'$   
**assumes** *R1*:  $i * k \leq j * k \implies R$   
**assumes** *R2*:  $\text{Suc } m * k' \leq j' * k' \implies R$   
**shows** *R*  $\langle \text{proof} \rangle$

**lemma** *lrlem*:  $(0 :: \text{nat}) < sc \implies$

$(sc - n + (n + lb * n) \leq m * n) = (sc + lb * n \leq m * n)$   
 $\langle proof \rangle$

**lemma** *gen-minus*:  $0 < n \implies f\ n = f\ (Suc\ (n - 1))$   
 $\langle proof \rangle$

**lemma** *mpl-lem*:  $j \leq (i :: nat) \implies k < j \implies i - j + k < i$   
 $\langle proof \rangle$

**lemma** *nonneg-mod-div*:  
 $0 \leq a \implies 0 \leq b \implies 0 \leq (a \bmod b :: int) \ \& \ 0 \leq a \div b$   
 $\langle proof \rangle$

**end**

## 4 BinGeneral: Basic Definitions for Binary Integers

**theory** *BinGeneral*  
**imports** *Num-Lemmas*  
**begin**

### 4.1 Further properties of numerals

**datatype** *bit* = *B0* | *B1*

**definition**  
 $Bit :: int \Rightarrow bit \Rightarrow int$  (**infixl** *BIT* 90) **where**  
 $k\ BIT\ b = (case\ b\ of\ B0 \Rightarrow 0 \mid B1 \Rightarrow 1) + k + k$

**lemma** *BIT-B0-eq-Bit0* [*simp*]:  $w\ BIT\ B0 = Int.Bit0\ w$   
 $\langle proof \rangle$

**lemma** *BIT-B1-eq-Bit1* [*simp*]:  $w\ BIT\ B1 = Int.Bit1\ w$   
 $\langle proof \rangle$

**lemmas** *BIT-simps* = *BIT-B0-eq-Bit0 BIT-B1-eq-Bit1*

**hide** (**open**) *const B0 B1*

**lemma** *Min-ne-Pls* [*iff*]:  
 $Int.Min \sim = Int.Pl$   
 $\langle proof \rangle$

**lemmas** *Pls-ne-Min* [*iff*] = *Min-ne-Pls* [*symmetric*]

**lemmas** *PlsMin-defs* [*intro!*] =



*Pls-def Min-def Pls-def [symmetric] Min-def [symmetric]*

**lemmas** *PlsMin-simps [simp] = PlsMin-defs [THEN Eq-TrueI]*

**lemma** *number-of-False-cong:*

*False  $\implies$  number-of  $x$  = number-of  $y$*   
 $\langle \text{proof} \rangle$

**lemma** *BIT-eq:  $u \text{ BIT } b = v \text{ BIT } c \implies u = v \ \& \ b = c$*   
 $\langle \text{proof} \rangle$

**lemmas** *BIT-eqE [elim!] = BIT-eq [THEN conjE, standard]*

**lemma** *BIT-eq-iff [simp]:*

*$(u \text{ BIT } b = v \text{ BIT } c) = (u = v \wedge b = c)$*   
 $\langle \text{proof} \rangle$

**lemmas** *BIT-eqI [intro!] = conjI [THEN BIT-eq-iff [THEN iffD2]]*

**lemma** *less-Bits:*

*$(v \text{ BIT } b < w \text{ BIT } c) = (v < w \mid v \leq w \ \& \ b = \text{bit.B0} \ \& \ c = \text{bit.B1})$*   
 $\langle \text{proof} \rangle$

**lemma** *le-Bits:*

*$(v \text{ BIT } b \leq w \text{ BIT } c) = (v < w \mid v \leq w \ \& \ (b \sim = \text{bit.B1} \mid c \sim = \text{bit.B0}))$*   
 $\langle \text{proof} \rangle$

**lemma** *no-no [simp]: number-of (number-of  $i$ ) =  $i$*   
 $\langle \text{proof} \rangle$

**lemma** *Bit-B0:*

*$k \text{ BIT } \text{bit.B0} = k + k$*   
 $\langle \text{proof} \rangle$

**lemma** *Bit-B1:*

*$k \text{ BIT } \text{bit.B1} = k + k + 1$*   
 $\langle \text{proof} \rangle$

**lemma** *Bit-B0-2t:  $k \text{ BIT } \text{bit.B0} = 2 * k$*   
 $\langle \text{proof} \rangle$

**lemma** *Bit-B1-2t:  $k \text{ BIT } \text{bit.B1} = 2 * k + 1$*   
 $\langle \text{proof} \rangle$

**lemma** *B-mod-2':*

*$X = 2 \implies (w \text{ BIT } \text{bit.B1}) \bmod X = 1 \ \& \ (w \text{ BIT } \text{bit.B0}) \bmod X = 0$*   
 $\langle \text{proof} \rangle$

**lemma** *B1-mod-2* [simp]:  $(\text{Int.Bit1 } w) \bmod 2 = 1$   
 ⟨proof⟩

**lemma** *B0-mod-2* [simp]:  $(\text{Int.Bit0 } w) \bmod 2 = 0$   
 ⟨proof⟩

**lemma** *neB1E* [elim!]:  
 assumes *ne*:  $y \neq \text{bit.B1}$   
 assumes *y*:  $y = \text{bit.B0} \implies P$   
 shows *P*  
 ⟨proof⟩

**lemma** *bin-ex-rl*:  $\text{EX } w \text{ b. } w \text{ BIT } b = \text{bin}$   
 ⟨proof⟩

**lemma** *bin-exhaust*:  
 assumes *Q*:  $\bigwedge x \text{ b. } \text{bin} = x \text{ BIT } b \implies Q$   
 shows *Q*  
 ⟨proof⟩

## 4.2 Destructors for binary integers

**definition** *bin-rl* ::  $\text{int} \Rightarrow \text{int} \times \text{bit}$  **where**  
 [code func del]:  $\text{bin-rl } w = (\text{THE } (r, l). w = r \text{ BIT } l)$

**lemma** *bin-rl-char*:  $(\text{bin-rl } w = (r, l)) = (r \text{ BIT } l = w)$   
 ⟨proof⟩

**definition**  
*bin-rest-def* [code func del]:  $\text{bin-rest } w = \text{fst } (\text{bin-rl } w)$

**definition**  
*bin-last-def* [code func del]:  $\text{bin-last } w = \text{snd } (\text{bin-rl } w)$

**primrec** *bin-nth* **where**  
*Z*:  $\text{bin-nth } w \ 0 = (\text{bin-last } w = \text{bit.B1})$   
 | *Suc*:  $\text{bin-nth } w \ (\text{Suc } n) = \text{bin-nth } (\text{bin-rest } w) \ n$

**lemma** *bin-rl*:  $\text{bin-rl } w = (\text{bin-rest } w, \text{bin-last } w)$   
 ⟨proof⟩

**lemma** *bin-rl-simps* [simp]:  
 $\text{bin-rl } \text{Int.Pls} = (\text{Int.Pls}, \text{bit.B0})$   
 $\text{bin-rl } \text{Int.Min} = (\text{Int.Min}, \text{bit.B1})$   
 $\text{bin-rl } (\text{Int.Bit0 } r) = (r, \text{bit.B0})$   
 $\text{bin-rl } (\text{Int.Bit1 } r) = (r, \text{bit.B1})$   
 $\text{bin-rl } (r \text{ BIT } b) = (r, b)$   
 ⟨proof⟩

**declare** *bin-rl-simps*(1-4) [code func]

**lemmas** *bin-rl-simp* [simp] = iffD1 [OF *bin-rl-char bin-rl*]

**lemma** *bin-abs-lem*:

*bin* = (*w BIT b*) ==> ~ *bin* = *Int.Min* --> ~ *bin* = *Int.Pl* -->  
 nat (*abs w*) < nat (*abs bin*)  
 <proof>

**lemma** *bin-induct*:

**assumes** *PPls*: *P Int.Pl*  
**and** *PMin*: *P Int.Min*  
**and** *PBit*: !!*bin bit. P bin* ==> *P (bin BIT bit)*  
**shows** *P bin*  
 <proof>

**lemma** *numeral-induct*:

**assumes** *Pls*: *P Int.Pl*  
**assumes** *Min*: *P Int.Min*  
**assumes** *Bit0*:  $\bigwedge w. \llbracket P\ w; w \neq \text{Int.Pl} \rrbracket \implies P\ (\text{Int.Bit0}\ w)$   
**assumes** *Bit1*:  $\bigwedge w. \llbracket P\ w; w \neq \text{Int.Min} \rrbracket \implies P\ (\text{Int.Bit1}\ w)$   
**shows** *P x*  
 <proof>

**lemma** *bin-rest-simps* [simp]:

*bin-rest Int.Pl* = *Int.Pl*  
*bin-rest Int.Min* = *Int.Min*  
*bin-rest (Int.Bit0 w)* = *w*  
*bin-rest (Int.Bit1 w)* = *w*  
*bin-rest (w BIT b)* = *w*  
 <proof>

**declare** *bin-rest-simps*(1-4) [code func]

**lemma** *bin-last-simps* [simp]:

*bin-last Int.Pl* = *bit.B0*  
*bin-last Int.Min* = *bit.B1*  
*bin-last (Int.Bit0 w)* = *bit.B0*  
*bin-last (Int.Bit1 w)* = *bit.B1*  
*bin-last (w BIT b)* = *b*  
 <proof>

**declare** *bin-last-simps*(1-4) [code func]

**lemma** *bin-r-l-extras* [simp]:

*bin-last 0* = *bit.B0*  
*bin-last (- 1)* = *bit.B1*  
*bin-last -1* = *bit.B1*

$bin\_last\ 1 = bit.B1$   
 $bin\_rest\ 1 = 0$   
 $bin\_rest\ 0 = 0$   
 $bin\_rest\ (-1) = -1$   
 $bin\_rest\ -1 = -1$   
 $\langle proof \rangle$

**lemma** *bin-last-mod*:

$bin\_last\ w = (if\ w\ mod\ 2 = 0\ then\ bit.B0\ else\ bit.B1)$   
 $\langle proof \rangle$

**lemma** *bin-rest-div*:

$bin\_rest\ w = w\ div\ 2$   
 $\langle proof \rangle$

**lemma** *Bit-div2 [simp]*:  $(w\ BIT\ b)\ div\ 2 = w$

$\langle proof \rangle$

**lemma** *Bit0-div2 [simp]*:  $(Int.Bit0\ w)\ div\ 2 = w$

$\langle proof \rangle$

**lemma** *Bit1-div2 [simp]*:  $(Int.Bit1\ w)\ div\ 2 = w$

$\langle proof \rangle$

**lemma** *bin-nth-lem [rule-format]*:

$ALL\ y.\ bin\_nth\ x = bin\_nth\ y \implies x = y$   
 $\langle proof \rangle$

**lemma** *bin-nth-eq-iff*:  $(bin\_nth\ x = bin\_nth\ y) = (x = y)$

$\langle proof \rangle$

**lemmas** *bin-eqI = ext [THEN bin-nth-eq-iff [THEN iffD1], standard]*

**lemma** *bin-nth-Pls [simp]*:  $\sim\ bin\_nth\ Int.Pl\ n$

$\langle proof \rangle$

**lemma** *bin-nth-Min [simp]*:  $bin\_nth\ Int.Min\ n$

$\langle proof \rangle$

**lemma** *bin-nth-0-BIT*:  $bin\_nth\ (w\ BIT\ b)\ 0 = (b = bit.B1)$

$\langle proof \rangle$

**lemma** *bin-nth-Suc-BIT*:  $bin\_nth\ (w\ BIT\ b)\ (Suc\ n) = bin\_nth\ w\ n$

$\langle proof \rangle$

**lemma** *bin-nth-minus [simp]*:  $0 < n \implies bin\_nth\ (w\ BIT\ b)\ n = bin\_nth\ w\ (n - 1)$

$\langle proof \rangle$

**lemma** *bin-nth-minus-Bit0* [simp]:  
 $0 < n \implies \text{bin-nth } (\text{Int.Bit0 } w) \ n = \text{bin-nth } w \ (n - 1)$   
 <proof>

**lemma** *bin-nth-minus-Bit1* [simp]:  
 $0 < n \implies \text{bin-nth } (\text{Int.Bit1 } w) \ n = \text{bin-nth } w \ (n - 1)$   
 <proof>

**lemmas** *bin-nth-0* = *bin-nth.simps*(1)  
**lemmas** *bin-nth-Suc* = *bin-nth.simps*(2)

**lemmas** *bin-nth-simps* =  
*bin-nth-0 bin-nth-Suc bin-nth-Pls bin-nth-Min bin-nth-minus*  
*bin-nth-minus-Bit0 bin-nth-minus-Bit1*

### 4.3 Recursion combinator for binary integers

**lemma** *brlem*:  $(\text{bin} = \text{Int.Min}) = (- \text{bin} + \text{Int.pred } 0 = 0)$   
 <proof>

**function**  
 $\text{bin-rec} :: 'a \Rightarrow 'a \Rightarrow (\text{int} \Rightarrow \text{bit} \Rightarrow 'a \Rightarrow 'a) \Rightarrow \text{int} \Rightarrow 'a$   
**where**  
 $\text{bin-rec } f1 \ f2 \ f3 \ \text{bin} = (\text{if } \text{bin} = \text{Int.Plus} \text{ then } f1$   
 $\quad \text{else if } \text{bin} = \text{Int.Min} \text{ then } f2$   
 $\quad \text{else case } \text{bin-rl } \text{bin} \text{ of } (w, b) \Rightarrow f3 \ w \ b \ (\text{bin-rec } f1 \ f2 \ f3 \ w))$   
 <proof>

**termination**  
 <proof>

**declare** *bin-rec.simps* [simp del]

**lemma** *bin-rec-PM*:  
 $f = \text{bin-rec } f1 \ f2 \ f3 \implies f \ \text{Int.Plus} = f1 \ \& \ f \ \text{Int.Min} = f2$   
 <proof>

**lemma** *bin-rec-Pls*:  $\text{bin-rec } f1 \ f2 \ f3 \ \text{Int.Plus} = f1$   
 <proof>

**lemma** *bin-rec-Min*:  $\text{bin-rec } f1 \ f2 \ f3 \ \text{Int.Min} = f2$   
 <proof>

**lemma** *bin-rec-Bit0*:  
 $f3 \ \text{Int.Plus} \ \text{bit.B0} \ f1 \implies$   
 $\text{bin-rec } f1 \ f2 \ f3 \ (\text{Int.Bit0 } w) = f3 \ w \ \text{bit.B0} \ (\text{bin-rec } f1 \ f2 \ f3 \ w)$   
 <proof>

**lemma** *bin-rec-Bit1*:

$f3 \text{ Int.Min } bit.B1 \ f2 = f2 \implies$   
 $\text{bin-rec } f1 \ f2 \ f3 \ (\text{Int.Bit1 } w) = f3 \ w \ bit.B1 \ (\text{bin-rec } f1 \ f2 \ f3 \ w)$   
 $\langle \text{proof} \rangle$

**lemma** *bin-rec-Bit*:

$f = \text{bin-rec } f1 \ f2 \ f3 \implies f3 \ \text{Int.Pls } bit.B0 \ f1 = f1 \implies$   
 $f3 \ \text{Int.Min } bit.B1 \ f2 = f2 \implies f \ (w \ \text{BIT } b) = f3 \ w \ b \ (f \ w)$   
 $\langle \text{proof} \rangle$

**lemmas** *bin-rec-simps* = *refl* [THEN *bin-rec-Bit*] *bin-rec-Pls bin-rec-Min bin-rec-Bit0 bin-rec-Bit1*

#### 4.4 Truncating binary integers

**definition**

*bin-sign-def* [code func del] : *bin-sign* = *bin-rec Int.Pls Int.Min* (%w b s. s)

**lemma** *bin-sign-simps* [simp]:

$\text{bin-sign } \text{Int.Pls} = \text{Int.Pls}$   
 $\text{bin-sign } \text{Int.Min} = \text{Int.Min}$   
 $\text{bin-sign } (\text{Int.Bit0 } w) = \text{bin-sign } w$   
 $\text{bin-sign } (\text{Int.Bit1 } w) = \text{bin-sign } w$   
 $\text{bin-sign } (w \ \text{BIT } b) = \text{bin-sign } w$   
 $\langle \text{proof} \rangle$

**declare** *bin-sign-simps*(1-4) [code func]

**lemma** *bin-sign-rest* [simp]:

$\text{bin-sign } (\text{bin-rest } w) = (\text{bin-sign } w)$   
 $\langle \text{proof} \rangle$

**consts**

*bintrunc* :: nat => int => int

**primrec**

$Z : \text{bintrunc } 0 \ \text{bin} = \text{Int.Pls}$   
 $\text{Suc} : \text{bintrunc } (\text{Suc } n) \ \text{bin} = \text{bintrunc } n \ (\text{bin-rest } \text{bin}) \ \text{BIT} \ (\text{bin-last } \text{bin})$

**consts**

*sbintrunc* :: nat => int => int

**primrec**

$Z : \text{sbintrunc } 0 \ \text{bin} =$   
 $(\text{case } \text{bin-last } \text{bin} \text{ of } bit.B1 \Rightarrow \text{Int.Min} \mid bit.B0 \Rightarrow \text{Int.Pls})$   
 $\text{Suc} : \text{sbintrunc } (\text{Suc } n) \ \text{bin} = \text{sbintrunc } n \ (\text{bin-rest } \text{bin}) \ \text{BIT} \ (\text{bin-last } \text{bin})$

**lemma** *sign-bintr*:

$!!w. \text{bin-sign } (\text{bintrunc } n \ w) = \text{Int.Pls}$   
 $\langle \text{proof} \rangle$

**lemma** *bintrunc-mod2p*:

$!!w. \text{bintrunc } n \ w = (w \bmod 2^n :: \text{int})$   
 $\langle \text{proof} \rangle$

**lemma** *sbintrunc-mod2p*:

$!!w. \text{sbintrunc } n \ w = ((w + 2^n) \bmod 2^{(\text{Suc } n) - 2^n} :: \text{int})$   
 $\langle \text{proof} \rangle$

#### 4.5 Simplifications for (s)bintrunc

**lemma** *bit-bool*:

$(b = (b' = \text{bit}.B1)) = (b' = (\text{if } b \text{ then } \text{bit}.B1 \text{ else } \text{bit}.B0))$   
 $\langle \text{proof} \rangle$

**lemmas** *bit-bool1* [simp] = refl [THEN bit-bool [THEN iffD1], symmetric]

**lemma** *bin-sign-lem*:

$!!\text{bin}. (\text{bin-sign } (\text{sbintrunc } n \ \text{bin}) = \text{Int.Min}) = \text{bin-nth } \text{bin } n$   
 $\langle \text{proof} \rangle$

**lemma** *nth-bintr*:

$!!w \ m. \text{bin-nth } (\text{bintrunc } m \ w) \ n = (n < m \ \& \ \text{bin-nth } w \ n)$   
 $\langle \text{proof} \rangle$

**lemma** *nth-sbintr*:

$!!w \ m. \text{bin-nth } (\text{sbintrunc } m \ w) \ n =$   
 $(\text{if } n < m \text{ then } \text{bin-nth } w \ n \text{ else } \text{bin-nth } w \ m)$   
 $\langle \text{proof} \rangle$

**lemma** *bin-nth-Bit*:

$\text{bin-nth } (w \ \text{BIT } b) \ n = (n = 0 \ \& \ b = \text{bit}.B1 \mid (\text{EX } m. n = \text{Suc } m \ \& \ \text{bin-nth } w \ m))$   
 $\langle \text{proof} \rangle$

**lemma** *bin-nth-Bit0*:

$\text{bin-nth } (\text{Int.Bit0 } w) \ n = (\text{EX } m. n = \text{Suc } m \ \& \ \text{bin-nth } w \ m)$   
 $\langle \text{proof} \rangle$

**lemma** *bin-nth-Bit1*:

$\text{bin-nth } (\text{Int.Bit1 } w) \ n = (n = 0 \mid (\text{EX } m. n = \text{Suc } m \ \& \ \text{bin-nth } w \ m))$   
 $\langle \text{proof} \rangle$

**lemma** *bintrunc-bintrunc-l*:

$n \leq m \implies (\text{bintrunc } m \ (\text{bintrunc } n \ w) = \text{bintrunc } n \ w)$   
 $\langle \text{proof} \rangle$

**lemma** *sbintrunc-sbintrunc-l*:

$n \leq m \implies (\text{sbintrunc } m \ (\text{sbintrunc } n \ w) = \text{sbintrunc } n \ w)$   
 $\langle \text{proof} \rangle$

**lemma** *bintrunc-bintrunc-ge*:

$$n \leq m \implies (\text{bintrunc } n (\text{bintrunc } m w) = \text{bintrunc } n w)$$

*<proof>*

**lemma** *bintrunc-bintrunc-min* [simp]:

$$\text{bintrunc } m (\text{bintrunc } n w) = \text{bintrunc } (\min m n) w$$

*<proof>*

**lemma** *sbintrunc-sbintrunc-min* [simp]:

$$\text{sbintrunc } m (\text{sbintrunc } n w) = \text{sbintrunc } (\min m n) w$$

*<proof>*

**lemmas** *bintrunc-Pls* =

*bintrunc.Suc [where bin=Int.Pls, simplified bin-last-simps bin-rest-simps, standard]*

**lemmas** *bintrunc-Min* [simp] =

*bintrunc.Suc [where bin=Int.Min, simplified bin-last-simps bin-rest-simps, standard]*

**lemmas** *bintrunc-BIT* [simp] =

*bintrunc.Suc [where bin=w BIT b, simplified bin-last-simps bin-rest-simps, standard]*

**lemma** *bintrunc-Bit0* [simp]:

$$\text{bintrunc } (\text{Suc } n) (\text{Int.Bit0 } w) = \text{Int.Bit0 } (\text{bintrunc } n w)$$

*<proof>*

**lemma** *bintrunc-Bit1* [simp]:

$$\text{bintrunc } (\text{Suc } n) (\text{Int.Bit1 } w) = \text{Int.Bit1 } (\text{bintrunc } n w)$$

*<proof>*

**lemmas** *bintrunc-Sucs* = *bintrunc-Pls bintrunc-Min bintrunc-BIT*

*bintrunc-Bit0 bintrunc-Bit1*

**lemmas** *sbintrunc-Suc-Pls* =

*sbintrunc.Suc [where bin=Int.Pls, simplified bin-last-simps bin-rest-simps, standard]*

**lemmas** *sbintrunc-Suc-Min* =

*sbintrunc.Suc [where bin=Int.Min, simplified bin-last-simps bin-rest-simps, standard]*

**lemmas** *sbintrunc-Suc-BIT* [simp] =

*sbintrunc.Suc [where bin=w BIT b, simplified bin-last-simps bin-rest-simps, standard]*

**lemma** *sbintrunc-Suc-Bit0* [simp]:

$$\text{sbintrunc } (\text{Suc } n) (\text{Int.Bit0 } w) = \text{Int.Bit0 } (\text{sbintrunc } n w)$$



$\langle \text{proof} \rangle$

**lemma** *sbintrunc-Suc-Bit1* [simp]:  
 $\text{sbintrunc } (\text{Suc } n) (\text{Int.Bit1 } w) = \text{Int.Bit1 } (\text{sbintrunc } n \ w)$   
 $\langle \text{proof} \rangle$

**lemmas** *sbintrunc-Sucs* = *sbintrunc-Suc-Pls* *sbintrunc-Suc-Min* *sbintrunc-Suc-BIT*  
*sbintrunc-Suc-Bit0* *sbintrunc-Suc-Bit1*

**lemmas** *sbintrunc-Pls* =  
 $\text{sbintrunc.Z } [\text{where } \text{bin} = \text{Int.Pl},$   
 $\text{simplified bin-last-simps bin-rest-simps bit.simps, standard}]$

**lemmas** *sbintrunc-Min* =  
 $\text{sbintrunc.Z } [\text{where } \text{bin} = \text{Int.Min},$   
 $\text{simplified bin-last-simps bin-rest-simps bit.simps, standard}]$

**lemmas** *sbintrunc-0-BIT-B0* [simp] =  
 $\text{sbintrunc.Z } [\text{where } \text{bin} = w \ \text{BIT } \text{bit.B0},$   
 $\text{simplified bin-last-simps bin-rest-simps bit.simps, standard}]$

**lemmas** *sbintrunc-0-BIT-B1* [simp] =  
 $\text{sbintrunc.Z } [\text{where } \text{bin} = w \ \text{BIT } \text{bit.B1},$   
 $\text{simplified bin-last-simps bin-rest-simps bit.simps, standard}]$

**lemma** *sbintrunc-0-Bit0* [simp]:  $\text{sbintrunc } 0 (\text{Int.Bit0 } w) = \text{Int.Pl}$   
 $\langle \text{proof} \rangle$

**lemma** *sbintrunc-0-Bit1* [simp]:  $\text{sbintrunc } 0 (\text{Int.Bit1 } w) = \text{Int.Min}$   
 $\langle \text{proof} \rangle$

**lemmas** *sbintrunc-0-simps* =  
*sbintrunc-Pls* *sbintrunc-Min* *sbintrunc-0-BIT-B0* *sbintrunc-0-BIT-B1*  
*sbintrunc-0-Bit0* *sbintrunc-0-Bit1*

**lemmas** *bintrunc-simps* = *bintrunc.Z* *bintrunc-Sucs*  
**lemmas** *sbintrunc-simps* = *sbintrunc-0-simps* *sbintrunc-Sucs*

**lemma** *bintrunc-minus*:  
 $0 < n ==> \text{bintrunc } (\text{Suc } (n - 1)) \ w = \text{bintrunc } n \ w$   
 $\langle \text{proof} \rangle$

**lemma** *sbintrunc-minus*:  
 $0 < n ==> \text{sbintrunc } (\text{Suc } (n - 1)) \ w = \text{sbintrunc } n \ w$   
 $\langle \text{proof} \rangle$

**lemmas** *bintrunc-minus-simps* =  
*bintrunc-Sucs* [THEN [2] *bintrunc-minus* [symmetric, THEN trans], standard]  
**lemmas** *sbintrunc-minus-simps* =

*sbintrunc-Sucs* [THEN [2] *sbintrunc-minus* [symmetric, THEN trans], standard]

**lemma** *bintrunc-n-Pls* [simp]:  
 $\text{bintrunc } n \text{ Int.Pls} = \text{Int.Pls}$   
 ⟨proof⟩

**lemma** *sbintrunc-n-PM* [simp]:  
 $\text{sbintrunc } n \text{ Int.Pls} = \text{Int.Pls}$   
 $\text{sbintrunc } n \text{ Int.Min} = \text{Int.Min}$   
 ⟨proof⟩

**lemmas** *thobini1* = arg-cong [where  $f = \%w. w \text{ BIT } b$ , standard]

**lemmas** *bintrunc-BIT-I* = trans [OF *bintrunc-BIT thobini1*]  
**lemmas** *bintrunc-Min-I* = trans [OF *bintrunc-Min thobini1*]

**lemmas** *bmsts* = *bintrunc-minus-simps*(1–3) [THEN *thobini1* [THEN [2] trans], standard]  
**lemmas** *bintrunc-Pls-minus-I* = *bmsts*(1)  
**lemmas** *bintrunc-Min-minus-I* = *bmsts*(2)  
**lemmas** *bintrunc-BIT-minus-I* = *bmsts*(3)

**lemma** *bintrunc-0-Min*:  $\text{bintrunc } 0 \text{ Int.Min} = \text{Int.Pls}$   
 ⟨proof⟩  
**lemma** *bintrunc-0-BIT*:  $\text{bintrunc } 0 (w \text{ BIT } b) = \text{Int.Pls}$   
 ⟨proof⟩

**lemma** *bintrunc-Suc-lem*:  
 $\text{bintrunc } (\text{Suc } n) x = y \implies m = \text{Suc } n \implies \text{bintrunc } m x = y$   
 ⟨proof⟩

**lemmas** *bintrunc-Suc-Ialts* =  
 $\text{bintrunc-Min-I}$  [THEN *bintrunc-Suc-lem*, standard]  
 $\text{bintrunc-BIT-I}$  [THEN *bintrunc-Suc-lem*, standard]

**lemmas** *sbintrunc-BIT-I* = trans [OF *sbintrunc-Suc-BIT thobini1*]

**lemmas** *sbintrunc-Suc-Is* =  
 $\text{sbintrunc-Sucs}(1-3)$  [THEN *thobini1* [THEN [2] trans], standard]

**lemmas** *sbintrunc-Suc-minus-Is* =  
 $\text{sbintrunc-minus-simps}(1-3)$  [THEN *thobini1* [THEN [2] trans], standard]

**lemma** *sbintrunc-Suc-lem*:  
 $\text{sbintrunc } (\text{Suc } n) x = y \implies m = \text{Suc } n \implies \text{sbintrunc } m x = y$   
 ⟨proof⟩

**lemmas** *sbintrunc-Suc-Ialts* =  
 $\text{sbintrunc-Suc-Is}$  [THEN *sbintrunc-Suc-lem*, standard]

**lemma** *sbintrunc-bintrunc-lt*:

$m > n \implies \text{sbintrunc } n \ (\text{bintrunc } m \ w) = \text{sbintrunc } n \ w$   
 $\langle \text{proof} \rangle$

**lemma** *bintrunc-sbintrunc-le*:

$m \leq \text{Suc } n \implies \text{bintrunc } m \ (\text{sbintrunc } n \ w) = \text{bintrunc } m \ w$   
 $\langle \text{proof} \rangle$

**lemmas** *bintrunc-sbintrunc [simp]* = *order-refl [THEN bintrunc-sbintrunc-le]*

**lemmas** *sbintrunc-bintrunc [simp]* = *lessI [THEN sbintrunc-bintrunc-lt]*

**lemmas** *bintrunc-bintrunc [simp]* = *order-refl [THEN bintrunc-bintrunc-l]*

**lemmas** *sbintrunc-sbintrunc [simp]* = *order-refl [THEN sbintrunc-sbintrunc-l]*

**lemma** *bintrunc-sbintrunc' [simp]*:

$0 < n \implies \text{bintrunc } n \ (\text{sbintrunc } (n - 1) \ w) = \text{bintrunc } n \ w$   
 $\langle \text{proof} \rangle$

**lemma** *sbintrunc-bintrunc' [simp]*:

$0 < n \implies \text{sbintrunc } (n - 1) \ (\text{bintrunc } n \ w) = \text{sbintrunc } (n - 1) \ w$   
 $\langle \text{proof} \rangle$

**lemma** *bin-sbin-eq-iff*:

$\text{bintrunc } (\text{Suc } n) \ x = \text{bintrunc } (\text{Suc } n) \ y \iff$   
 $\text{sbintrunc } n \ x = \text{sbintrunc } n \ y$   
 $\langle \text{proof} \rangle$

**lemma** *bin-sbin-eq-iff'*:

$0 < n \implies \text{bintrunc } n \ x = \text{bintrunc } n \ y \iff$   
 $\text{sbintrunc } (n - 1) \ x = \text{sbintrunc } (n - 1) \ y$   
 $\langle \text{proof} \rangle$

**lemmas** *bintrunc-sbintruncS0 [simp]* = *bintrunc-sbintrunc' [unfolded One-nat-def]*

**lemmas** *sbintrunc-bintruncS0 [simp]* = *sbintrunc-bintrunc' [unfolded One-nat-def]*

**lemmas** *bintrunc-bintrunc-l' = le-add1 [THEN bintrunc-bintrunc-l]*

**lemmas** *sbintrunc-sbintrunc-l' = le-add1 [THEN sbintrunc-sbintrunc-l]*

**lemmas** *nat-non0-gr =*

*trans [OF iszero-def [THEN Not-eq-iff [THEN iffD2]] refl, standard]*

**lemmas** *bintrunc-pred-simps [simp] =*

*bintrunc-minus-simps [of number-of bin, simplified nobm1, standard]*

**lemmas** *sbintrunc-pred-simps [simp] =*

*sbintrunc-minus-simps [of number-of bin, simplified nobm1, standard]*

**lemma** *no-bintr-alt*:

$$\text{number-of } (\text{bintrunc } n \ w) = w \bmod 2^n$$

*<proof>*

**lemma** *no-bintr-alt1*:  $\text{bintrunc } n = (\%w. w \bmod 2^n :: \text{int})$

*<proof>*

**lemma** *range-bintrunc*:  $\text{range } (\text{bintrunc } n) = \{i. 0 \leq i \ \& \ i < 2^n\}$

*<proof>*

**lemma** *no-bintr*:

$$\text{number-of } (\text{bintrunc } n \ w) = (\text{number-of } w \bmod 2^n :: \text{int})$$

*<proof>*

**lemma** *no-sbintr-alt2*:

$$\text{sbintrunc } n = (\%w. (w + 2^n) \bmod 2^{\text{Suc } n} - 2^n :: \text{int})$$

*<proof>*

**lemma** *no-sbintr*:

$$\begin{aligned} \text{number-of } (\text{sbintrunc } n \ w) = \\ ((\text{number-of } w + 2^n) \bmod 2^{\text{Suc } n} - 2^n :: \text{int}) \end{aligned}$$

*<proof>*

**lemma** *range-sbintrunc*:

$$\text{range } (\text{sbintrunc } n) = \{i. -(2^n) \leq i \ \& \ i < 2^n\}$$

*<proof>*

**lemma** *sb-inc-lem*:

$$(a::\text{int}) + 2^k < 0 \implies a + 2^k + 2^{\text{Suc } k} \leq (a + 2^k) \bmod 2^{\text{Suc } k}$$

*<proof>*

**lemma** *sb-inc-lem'*:

$$(a::\text{int}) < -(2^k) \implies a + 2^k + 2^{\text{Suc } k} \leq (a + 2^k) \bmod 2^{\text{Suc } k}$$

*<proof>*

**lemma** *sbintrunc-inc*:

$$x < -(2^n) \implies x + 2^{\text{Suc } n} \leq \text{sbintrunc } n \ x$$

*<proof>*

**lemma** *sb-dec-lem*:

$$(0::\text{int}) \leq -(2^k) + a \implies (a + 2^k) \bmod (2 * 2^k) \leq -(2^k) + a$$

*<proof>*

**lemma** *sb-dec-lem'*:

$$(2::\text{int})^k \leq a \implies (a + 2^k) \bmod (2 * 2^k) \leq -(2^k) + a$$

*<proof>*

**lemma** *sbintrunc-dec*:

$$x \geq (2^n) \implies x - 2^{\text{Suc } n} \geq \text{sbintrunc } n \ x$$

$\langle \text{proof} \rangle$

**lemmas**  $\text{zmod-uminus}' = \text{zmod-uminus}$  [where  $b=c$ , standard]

**lemmas**  $\text{zpower-zmod}' = \text{zpower-zmod}$  [where  $m=c$  and  $y=k$ , standard]

**lemmas**  $\text{brdmod1s}'$  [symmetric] =  
 $\text{zmod-zadd-left-eq}$   $\text{zmod-zadd-right-eq}$   
 $\text{zmod-zsub-left-eq}$   $\text{zmod-zsub-right-eq}$   
 $\text{zmod-zmult1-eq}$   $\text{zmod-zmult1-eq-rev}$

**lemmas**  $\text{brdmods}'$  [symmetric] =  
 $\text{zpower-zmod}'$  [symmetric]  
 $\text{trans}$  [OF  $\text{zmod-zadd-left-eq}$   $\text{zmod-zadd-right-eq}$ ]  
 $\text{trans}$  [OF  $\text{zmod-zsub-left-eq}$   $\text{zmod-zsub-right-eq}$ ]  
 $\text{trans}$  [OF  $\text{zmod-zmult1-eq}$   $\text{zmod-zmult1-eq-rev}$ ]  
 $\text{zmod-uminus}'$  [symmetric]  
 $\text{zmod-zadd-left-eq}$  [where  $b = 1$ ]  
 $\text{zmod-zsub-left-eq}$  [where  $b = 1$ ]

**lemmas**  $\text{bintr-arith1s} =$   
 $\text{brdmod1s}'$  [where  $c=2^n$ , folded pred-def succ-def bintrunc-mod2p, standard]

**lemmas**  $\text{bintr-ariths} =$   
 $\text{brdmods}'$  [where  $c=2^n$ , folded pred-def succ-def bintrunc-mod2p, standard]

**lemmas**  $\text{m2pths} = \text{pos-mod-sign pos-mod-bound}$  [OF  $\text{zless2p}$ , standard]

**lemma**  $\text{bintr-ge0}$ :  $(0 :: \text{int}) \leq \text{number-of } (\text{bintrunc } n \ w)$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bintr-lt2p}$ :  $\text{number-of } (\text{bintrunc } n \ w) < (2^n :: \text{int})$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bintr-Min}$ :  
 $\text{number-of } (\text{bintrunc } n \ \text{Int.Min}) = (2^n :: \text{int}) - 1$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{sbintr-ge}$ :  $(-(2^n :: \text{int})) \leq \text{number-of } (\text{sbintrunc } n \ w)$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{sbintr-lt}$ :  $\text{number-of } (\text{sbintrunc } n \ w) < (2^n :: \text{int})$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bintrunc-Suc}$ :  
 $\text{bintrunc } (\text{Suc } n) \ \text{bin} = \text{bintrunc } n \ (\text{bin-rest bin}) \ \text{BIT bin-last bin}$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{sign-Pls-ge-0}$ :  
 $(\text{bin-sign bin} = \text{Int.Pls}) = (\text{number-of bin} \geq (0 :: \text{int}))$   
 $\langle \text{proof} \rangle$

**lemma** *sign-Min-lt-0*:

$(\text{bin-sign bin} = \text{Int.Min}) = (\text{number-of bin} < (0 :: \text{int}))$   
 $\langle \text{proof} \rangle$

**lemmas** *sign-Min-neg* = *trans* [*OF sign-Min-lt-0 neg-def* [*symmetric*]]

**lemma** *bin-rest-trunc*:

$!!\text{bin. } (\text{bin-rest } (\text{bintrunc } n \text{ bin})) = \text{bintrunc } (n - 1) (\text{bin-rest bin})$   
 $\langle \text{proof} \rangle$

**lemma** *bin-rest-power-trunc* [*rule-format*] :

$(\text{bin-rest } ^k) (\text{bintrunc } n \text{ bin}) =$   
 $\text{bintrunc } (n - k) ((\text{bin-rest } ^k) \text{ bin})$   
 $\langle \text{proof} \rangle$

**lemma** *bin-rest-trunc-i*:

$\text{bintrunc } n (\text{bin-rest bin}) = \text{bin-rest } (\text{bintrunc } (\text{Suc } n) \text{ bin})$   
 $\langle \text{proof} \rangle$

**lemma** *bin-rest-strunc*:

$!!\text{bin. } \text{bin-rest } (\text{sbintrunc } (\text{Suc } n) \text{ bin}) = \text{sbintrunc } n (\text{bin-rest bin})$   
 $\langle \text{proof} \rangle$

**lemma** *bintrunc-rest* [*simp*]:

$!!\text{bin. } \text{bintrunc } n (\text{bin-rest } (\text{bintrunc } n \text{ bin})) = \text{bin-rest } (\text{bintrunc } n \text{ bin})$   
 $\langle \text{proof} \rangle$

**lemma** *sbintrunc-rest* [*simp*]:

$!!\text{bin. } \text{sbintrunc } n (\text{bin-rest } (\text{sbintrunc } n \text{ bin})) = \text{bin-rest } (\text{sbintrunc } n \text{ bin})$   
 $\langle \text{proof} \rangle$

**lemma** *bintrunc-rest'*:

$\text{bintrunc } n \circ \text{bin-rest} \circ \text{bintrunc } n = \text{bin-rest} \circ \text{bintrunc } n$   
 $\langle \text{proof} \rangle$

**lemma** *sbintrunc-rest'* :

$\text{sbintrunc } n \circ \text{bin-rest} \circ \text{sbintrunc } n = \text{bin-rest} \circ \text{sbintrunc } n$   
 $\langle \text{proof} \rangle$

**lemma** *rco-lem*:

$f \circ g \circ f = g \circ f \implies f \circ (g \circ f) ^n = g ^n \circ f$   
 $\langle \text{proof} \rangle$

**lemma** *rco-alt*:  $(f \circ g) ^n \circ f = f \circ (g \circ f) ^n$

$\langle \text{proof} \rangle$

**lemmas** *rco-bintr* = *bintrunc-rest'*

[*THEN rco-lem* [*THEN fun-cong*], *unfolded o-def*]

**lemmas** *rco-sbintr* = *sbintrunc-rest'*  
 [THEN *rco-lem* [THEN *fun-cong*], unfolded *o-def*]

#### 4.6 Splitting and concatenation

**primrec** *bin-split* :: *nat*  $\Rightarrow$  *int*  $\Rightarrow$  *int*  $\times$  *int* **where**  
*Z*: *bin-split* 0 *w* = (*w*, *Int.Pls*)  
 | *Suc*: *bin-split* (*Suc* *n*) *w* = (*let* (*w1*, *w2*) = *bin-split* *n* (*bin-rest* *w*)  
   in (*w1*, *w2* BIT *bin-last* *w*))

**primrec** *bin-cat* :: *int*  $\Rightarrow$  *nat*  $\Rightarrow$  *int*  $\Rightarrow$  *int* **where**  
*Z*: *bin-cat* *w* 0 *v* = *w*  
 | *Suc*: *bin-cat* *w* (*Suc* *n*) *v* = *bin-cat* *w* *n* (*bin-rest* *v*) BIT *bin-last* *v*

#### 4.7 Miscellaneous lemmas

**lemmas** *funpow-minus-simp* =  
*trans* [OF *gen-minus* [where *f* = *power* *f*] *funpow-Suc*, *standard*]

**lemmas** *funpow-pred-simp* [*simp*] =  
*funpow-minus-simp* [of *number-of* *bin*, *simplified nobm1*, *standard*]

**lemmas** *replicate-minus-simp* =  
*trans* [OF *gen-minus* [where *f* = %*n*. *replicate* *n* *x*] *replicate.replicate-Suc*,  
*standard*]

**lemmas** *replicate-pred-simp* [*simp*] =  
*replicate-minus-simp* [of *number-of* *bin*, *simplified nobm1*, *standard*]

**lemmas** *power-Suc-no* [*simp*] = *power-Suc* [of *number-of* *a*, *standard*]

**lemmas** *power-minus-simp* =  
*trans* [OF *gen-minus* [where *f* = *power* *f*] *power-Suc*, *standard*]

**lemmas** *power-pred-simp* =  
*power-minus-simp* [of *number-of* *bin*, *simplified nobm1*, *standard*]

**lemmas** *power-pred-simp-no* [*simp*] = *power-pred-simp* [where *f* = *number-of* *f*,  
*standard*]

**lemma** *list-exhaust-size-gt0*:  
*assumes* *y*:  $\bigwedge a \text{ list. } y = a \# \text{ list} \Rightarrow P$   
*shows*  $0 < \text{length } y \Rightarrow P$   
 <proof>

**lemma** *list-exhaust-size-eq0*:  
*assumes* *y*:  $y = [] \Rightarrow P$   
*shows*  $\text{length } y = 0 \Rightarrow P$   
 <proof>

**lemma** *size-Cons-lem-eq*:

$y = xa \# list \implies size\ y = Suc\ k \implies size\ list = k$   
 ⟨proof⟩

**lemma** *size-Cons-lem-eq-bin*:

$y = xa \# list \implies size\ y = number-of\ (Int.succ\ k) \implies$   
 $size\ list = number-of\ k$   
 ⟨proof⟩

**lemmas** *ls-splits* =

*prod.split split-split prod.split-asm split-split-asm split-if-asm*

**lemma** *not-B1-is-B0*:  $y \neq bit.B1 \implies y = bit.B0$

⟨proof⟩

**lemma** *B1-ass-B0*:

**assumes**  $y: y = bit.B0 \implies y = bit.B1$

**shows**  $y = bit.B1$

⟨proof⟩

**lemmas** *n2s-ths* [*THEN eq-reflection*] = *add-2-eq-Suc add-2-eq-Suc'*

**lemmas** *s2n-ths* = *n2s-ths* [*symmetric*]

**end**

## 5 BitSyntax: Syntactic classes for bitwise operations

**theory** *BitSyntax*

**imports** *BinGeneral*

**begin**

**class** *bit* = *type* +

**fixes** *bitNOT* ::  $'a \Rightarrow 'a$  (*NOT* - [70] 71)  
**and** *bitAND* ::  $'a \Rightarrow 'a \Rightarrow 'a$  (**infixr** *AND* 64)  
**and** *bitOR* ::  $'a \Rightarrow 'a \Rightarrow 'a$  (**infixr** *OR* 59)  
**and** *bitXOR* ::  $'a \Rightarrow 'a \Rightarrow 'a$  (**infixr** *XOR* 59)

We want the bitwise operations to bind slightly weaker than + and −, but ~ to bind slightly stronger than \*.

Testing and shifting operations.

**class** *bits* = *bit* +

**fixes** *test-bit* ::  $'a \Rightarrow nat \Rightarrow bool$  (**infixl** !! 100)  
**and** *lsb* ::  $'a \Rightarrow bool$   
**and** *set-bit* ::  $'a \Rightarrow nat \Rightarrow bool \Rightarrow 'a$   
**and** *set-bits* ::  $(nat \Rightarrow bool) \Rightarrow 'a$  (**binder** *BITS* 10)  
**and** *shiffl* ::  $'a \Rightarrow nat \Rightarrow 'a$  (**infixl** << 55)



**and** *shiftr* :: 'a  $\Rightarrow$  nat  $\Rightarrow$  'a (**infixl** >> 55)

**class** *bitss* = *bits* +  
**fixes** *msb* :: 'a  $\Rightarrow$  bool

### 5.1 Bitwise operations on *bit*

**instantiation** *bit* :: *bit*  
**begin**

**primrec** *bitNOT-bit* **where**  
 $NOT\ bit.B0 = bit.B1$   
 $| NOT\ bit.B1 = bit.B0$

**primrec** *bitAND-bit* **where**  
 $bit.B0\ AND\ y = bit.B0$   
 $| bit.B1\ AND\ y = y$

**primrec** *bitOR-bit* **where**  
 $bit.B0\ OR\ y = y$   
 $| bit.B1\ OR\ y = bit.B1$

**primrec** *bitXOR-bit* **where**  
 $bit.B0\ XOR\ y = y$   
 $| bit.B1\ XOR\ y = NOT\ y$

**instance**  $\langle proof \rangle$

**end**

**lemmas** *bit-simps* =  
 $bitNOT-bit.simps\ bitAND-bit.simps\ bitOR-bit.simps\ bitXOR-bit.simps$

**lemma** *bit-extra-simps* [*simp*]:  
 $x\ AND\ bit.B0 = bit.B0$   
 $x\ AND\ bit.B1 = x$   
 $x\ OR\ bit.B1 = bit.B1$   
 $x\ OR\ bit.B0 = x$   
 $x\ XOR\ bit.B1 = NOT\ x$   
 $x\ XOR\ bit.B0 = x$   
 $\langle proof \rangle$

**lemma** *bit-ops-comm*:  
 $(x::bit)\ AND\ y = y\ AND\ x$   
 $(x::bit)\ OR\ y = y\ OR\ x$   
 $(x::bit)\ XOR\ y = y\ XOR\ x$   
 $\langle proof \rangle$

**lemma** *bit-ops-same* [*simp*]:

```

(x::bit) AND x = x
(x::bit) OR x = x
(x::bit) XOR x = bit.B0
⟨proof⟩

```

```

lemma bit-not-not [simp]: NOT (NOT (x::bit)) = x
⟨proof⟩

```

```

end

```

## 6 BinOperations: Bitwise Operations on Binary Integers

```

theory BinOperations
imports BinGeneral BitSyntax
begin

```

### 6.1 Logical operations

bit-wise logical operations on the int type

```

instantiation int :: bit
begin

```

**definition**

```

int-not-def [code func del]: bitNOT = bin-rec Int.Min Int.Pls
  (λw b s. s BIT (NOT b))

```

**definition**

```

int-and-def [code func del]: bitAND = bin-rec (λx. Int.Pls) (λy. y)
  (λw b s y. s (bin-rest y) BIT (b AND bin-last y))

```

**definition**

```

int-or-def [code func del]: bitOR = bin-rec (λx. x) (λy. Int.Min)
  (λw b s y. s (bin-rest y) BIT (b OR bin-last y))

```

**definition**

```

int-xor-def [code func del]: bitXOR = bin-rec (λx. x) bitNOT
  (λw b s y. s (bin-rest y) BIT (b XOR bin-last y))

```

```

instance ⟨proof⟩

```

```

end

```

```

lemma int-not-simps [simp]:
  NOT Int.Pls = Int.Min
  NOT Int.Min = Int.Pls
  NOT (Int.Bit0 w) = Int.Bit1 (NOT w)

```

$NOT (Int.Bit1\ w) = Int.Bit0\ (NOT\ w)$   
 $NOT\ (w\ BIT\ b) = (NOT\ w)\ BIT\ (NOT\ b)$   
 $\langle proof \rangle$

**declare** *int-not-simps*(1-4) [code func]

**lemma** *int-xor-Pls* [simp, code func]:  
 $Int.Pl\ XOR\ x = x$   
 $\langle proof \rangle$

**lemma** *int-xor-Min* [simp, code func]:  
 $Int.Min\ XOR\ x = NOT\ x$   
 $\langle proof \rangle$

**lemma** *int-xor-Bits* [simp]:  
 $(x\ BIT\ b)\ XOR\ (y\ BIT\ c) = (x\ XOR\ y)\ BIT\ (b\ XOR\ c)$   
 $\langle proof \rangle$

**lemma** *int-xor-Bits2* [simp, code func]:  
 $(Int.Bit0\ x)\ XOR\ (Int.Bit0\ y) = Int.Bit0\ (x\ XOR\ y)$   
 $(Int.Bit0\ x)\ XOR\ (Int.Bit1\ y) = Int.Bit1\ (x\ XOR\ y)$   
 $(Int.Bit1\ x)\ XOR\ (Int.Bit0\ y) = Int.Bit1\ (x\ XOR\ y)$   
 $(Int.Bit1\ x)\ XOR\ (Int.Bit1\ y) = Int.Bit0\ (x\ XOR\ y)$   
 $\langle proof \rangle$

**lemma** *int-xor-x-simps'*:  
 $w\ XOR\ (Int.Pl\ BIT\ bit.B0) = w$   
 $w\ XOR\ (Int.Min\ BIT\ bit.B1) = NOT\ w$   
 $\langle proof \rangle$

**lemma** *int-xor-extra-simps* [simp, code func]:  
 $w\ XOR\ Int.Pl = w$   
 $w\ XOR\ Int.Min = NOT\ w$   
 $\langle proof \rangle$

**lemma** *int-or-Pls* [simp, code func]:  
 $Int.Pl\ OR\ x = x$   
 $\langle proof \rangle$

**lemma** *int-or-Min* [simp, code func]:  
 $Int.Min\ OR\ x = Int.Min$   
 $\langle proof \rangle$

**lemma** *int-or-Bits* [simp]:  
 $(x\ BIT\ b)\ OR\ (y\ BIT\ c) = (x\ OR\ y)\ BIT\ (b\ OR\ c)$   
 $\langle proof \rangle$

**lemma** *int-or-Bits2* [simp, code func]:  
 $(Int.Bit0\ x)\ OR\ (Int.Bit0\ y) = Int.Bit0\ (x\ OR\ y)$

$(Int.Bit0\ x)\ OR\ (Int.Bit1\ y) = Int.Bit1\ (x\ OR\ y)$   
 $(Int.Bit1\ x)\ OR\ (Int.Bit0\ y) = Int.Bit1\ (x\ OR\ y)$   
 $(Int.Bit1\ x)\ OR\ (Int.Bit1\ y) = Int.Bit1\ (x\ OR\ y)$   
 $\langle proof \rangle$

**lemma** *int-or-x-simps'*:

$w\ OR\ (Int.Pls\ BIT\ bit.B0) = w$   
 $w\ OR\ (Int.Min\ BIT\ bit.B1) = Int.Min$   
 $\langle proof \rangle$

**lemma** *int-or-extra-simps* [*simp*, *code func*]:

$w\ OR\ Int.Pls = w$   
 $w\ OR\ Int.Min = Int.Min$   
 $\langle proof \rangle$

**lemma** *int-and-Pls* [*simp*, *code func*]:

$Int.Pls\ AND\ x = Int.Pls$   
 $\langle proof \rangle$

**lemma** *int-and-Min* [*simp*, *code func*]:

$Int.Min\ AND\ x = x$   
 $\langle proof \rangle$

**lemma** *int-and-Bits* [*simp*]:

$(x\ BIT\ b)\ AND\ (y\ BIT\ c) = (x\ AND\ y)\ BIT\ (b\ AND\ c)$   
 $\langle proof \rangle$

**lemma** *int-and-Bits2* [*simp*, *code func*]:

$(Int.Bit0\ x)\ AND\ (Int.Bit0\ y) = Int.Bit0\ (x\ AND\ y)$   
 $(Int.Bit0\ x)\ AND\ (Int.Bit1\ y) = Int.Bit0\ (x\ AND\ y)$   
 $(Int.Bit1\ x)\ AND\ (Int.Bit0\ y) = Int.Bit0\ (x\ AND\ y)$   
 $(Int.Bit1\ x)\ AND\ (Int.Bit1\ y) = Int.Bit1\ (x\ AND\ y)$   
 $\langle proof \rangle$

**lemma** *int-and-x-simps'*:

$w\ AND\ (Int.Pls\ BIT\ bit.B0) = Int.Pls$   
 $w\ AND\ (Int.Min\ BIT\ bit.B1) = w$   
 $\langle proof \rangle$

**lemma** *int-and-extra-simps* [*simp*, *code func*]:

$w\ AND\ Int.Pls = Int.Pls$   
 $w\ AND\ Int.Min = w$   
 $\langle proof \rangle$

**lemma** *bin-ops-comm*:

**shows**

*int-and-comm*:  $!!y::int.\ x\ AND\ y = y\ AND\ x$  **and**  
*int-or-comm*:  $!!y::int.\ x\ OR\ y = y\ OR\ x$  **and**

*int-xor-comm*:  $!!y::int. x \text{ XOR } y = y \text{ XOR } x$   
 $\langle \text{proof} \rangle$

**lemma** *bin-ops-same* [simp]:

$(x::int) \text{ AND } x = x$   
 $(x::int) \text{ OR } x = x$   
 $(x::int) \text{ XOR } x = \text{Int.Pls}$   
 $\langle \text{proof} \rangle$

**lemma** *int-not-not* [simp]:  $\text{NOT } (\text{NOT } (x::int)) = x$   
 $\langle \text{proof} \rangle$

**lemmas** *bin-log-esimps* =

*int-and-extra-simps int-or-extra-simps int-xor-extra-simps*  
*int-and-Pls int-and-Min int-or-Pls int-or-Min int-xor-Pls int-xor-Min*

**lemma** *bbw-ao-absorb*:

$!!y::int. x \text{ AND } (y \text{ OR } x) = x \ \& \ x \text{ OR } (y \text{ AND } x) = x$   
 $\langle \text{proof} \rangle$

**lemma** *bbw-ao-absorbs-other*:

$x \text{ AND } (x \text{ OR } y) = x \wedge (y \text{ AND } x) \text{ OR } x = (x::int)$   
 $(y \text{ OR } x) \text{ AND } x = x \wedge x \text{ OR } (x \text{ AND } y) = (x::int)$   
 $(x \text{ OR } y) \text{ AND } x = x \wedge (x \text{ AND } y) \text{ OR } x = (x::int)$   
 $\langle \text{proof} \rangle$

**lemmas** *bbw-ao-absorbs* [simp] = *bbw-ao-absorb bbw-ao-absorbs-other*

**lemma** *int-xor-not*:

$!!y::int. (\text{NOT } x) \text{ XOR } y = \text{NOT } (x \text{ XOR } y) \ \&$   
 $x \text{ XOR } (\text{NOT } y) = \text{NOT } (x \text{ XOR } y)$   
 $\langle \text{proof} \rangle$

**lemma** *bbw-assocs'*:

$!!y \ z::int. (x \text{ AND } y) \text{ AND } z = x \text{ AND } (y \text{ AND } z) \ \&$   
 $(x \text{ OR } y) \text{ OR } z = x \text{ OR } (y \text{ OR } z) \ \&$   
 $(x \text{ XOR } y) \text{ XOR } z = x \text{ XOR } (y \text{ XOR } z)$   
 $\langle \text{proof} \rangle$

**lemma** *int-and-assoc*:

$(x \text{ AND } y) \text{ AND } (z::int) = x \text{ AND } (y \text{ AND } z)$   
 $\langle \text{proof} \rangle$

**lemma** *int-or-assoc*:

$(x \text{ OR } y) \text{ OR } (z::int) = x \text{ OR } (y \text{ OR } z)$   
 $\langle \text{proof} \rangle$

**lemma** *int-xor-assoc*:

$(x \text{ XOR } y) \text{ XOR } (z::\text{int}) = x \text{ XOR } (y \text{ XOR } z)$   
 $\langle \text{proof} \rangle$

**lemmas** *bbw-assocs* = *int-and-assoc int-or-assoc int-xor-assoc*

**lemma** *bbw-lcs [simp]*:

$(y::\text{int}) \text{ AND } (x \text{ AND } z) = x \text{ AND } (y \text{ AND } z)$   
 $(y::\text{int}) \text{ OR } (x \text{ OR } z) = x \text{ OR } (y \text{ OR } z)$   
 $(y::\text{int}) \text{ XOR } (x \text{ XOR } z) = x \text{ XOR } (y \text{ XOR } z)$   
 $\langle \text{proof} \rangle$

**lemma** *bbw-not-dist*:

$!!y::\text{int}. \text{ NOT } (x \text{ OR } y) = (\text{ NOT } x) \text{ AND } (\text{ NOT } y)$   
 $!!y::\text{int}. \text{ NOT } (x \text{ AND } y) = (\text{ NOT } x) \text{ OR } (\text{ NOT } y)$   
 $\langle \text{proof} \rangle$

**lemma** *bbw-oa-dist*:

$!!y \ z::\text{int}. (x \text{ AND } y) \text{ OR } z =$   
 $(x \text{ OR } z) \text{ AND } (y \text{ OR } z)$   
 $\langle \text{proof} \rangle$

**lemma** *bbw-ao-dist*:

$!!y \ z::\text{int}. (x \text{ OR } y) \text{ AND } z =$   
 $(x \text{ AND } z) \text{ OR } (y \text{ AND } z)$   
 $\langle \text{proof} \rangle$

**lemma** *plus-and-or [rule-format]*:

$\text{ ALL } y::\text{int}. (x \text{ AND } y) + (x \text{ OR } y) = x + y$   
 $\langle \text{proof} \rangle$

**lemma** *le-int-or*:

$!!x. \text{ bin-sign } y = \text{ Int.Pls } ==> x <= x \text{ OR } y$   
 $\langle \text{proof} \rangle$

**lemmas** *int-and-le* =

*xtr3 [OF bbw-ao-absorbs (2) [THEN conjunct2, symmetric] le-int-or]*

**lemma** *bin-nth-ops*:

$!!x \ y. \text{ bin-nth } (x \text{ AND } y) \ n = (\text{ bin-nth } x \ n \ \& \ \text{ bin-nth } y \ n)$   
 $!!x \ y. \text{ bin-nth } (x \text{ OR } y) \ n = (\text{ bin-nth } x \ n \ | \ \text{ bin-nth } y \ n)$   
 $!!x \ y. \text{ bin-nth } (x \text{ XOR } y) \ n = (\text{ bin-nth } x \ n \ \sim = \text{ bin-nth } y \ n)$   
 $!!x. \text{ bin-nth } (\text{ NOT } x) \ n = (\sim \text{ bin-nth } x \ n)$   
 $\langle \text{proof} \rangle$

**lemma** *bin-add-not*:  $x + \text{NOT } x = \text{Int.Min}$   
 ⟨proof⟩

**lemma** *bin-trunc-ao*:  
 $\text{!!}x\ y. (\text{bintrunc } n\ x) \text{ AND } (\text{bintrunc } n\ y) = \text{bintrunc } n\ (x \text{ AND } y)$   
 $\text{!!}x\ y. (\text{bintrunc } n\ x) \text{ OR } (\text{bintrunc } n\ y) = \text{bintrunc } n\ (x \text{ OR } y)$   
 ⟨proof⟩

**lemma** *bin-trunc-xor*:  
 $\text{!!}x\ y. \text{bintrunc } n\ (\text{bintrunc } n\ x \text{ XOR } \text{bintrunc } n\ y) =$   
 $\text{bintrunc } n\ (x \text{ XOR } y)$   
 ⟨proof⟩

**lemma** *bin-trunc-not*:  
 $\text{!!}x. \text{bintrunc } n\ (\text{NOT } (\text{bintrunc } n\ x)) = \text{bintrunc } n\ (\text{NOT } x)$   
 ⟨proof⟩

**lemma** *bintr-bintr-i*:  
 $x = \text{bintrunc } n\ y \implies \text{bintrunc } n\ x = \text{bintrunc } n\ y$   
 ⟨proof⟩

**lemmas** *bin-trunc-and* = *bin-trunc-ao*(1) [THEN *bintr-bintr-i*]  
**lemmas** *bin-trunc-or* = *bin-trunc-ao*(2) [THEN *bintr-bintr-i*]

## 6.2 Setting and clearing bits

**primrec**  
 $\text{bin-sc} :: \text{nat} \Rightarrow \text{bit} \Rightarrow \text{int} \Rightarrow \text{int}$   
**where**  
 $Z: \text{bin-sc } 0\ b\ w = \text{bin-rest } w\ \text{BIT } b$   
 $| \text{Suc}: \text{bin-sc } (\text{Suc } n)\ b\ w = \text{bin-sc } n\ b\ (\text{bin-rest } w)\ \text{BIT } \text{bin-last } w$

**lemma** *bin-nth-sc* [simp]:  
 $\text{!!}w. \text{bin-nth } (\text{bin-sc } n\ b\ w)\ n = (b = \text{bit.B1})$   
 ⟨proof⟩

**lemma** *bin-sc-sc-same* [simp]:  
 $\text{!!}w. \text{bin-sc } n\ c\ (\text{bin-sc } n\ b\ w) = \text{bin-sc } n\ c\ w$   
 ⟨proof⟩

**lemma** *bin-sc-sc-diff*:  
 $\text{!!}w\ m. m \sim n \implies$   
 $\text{bin-sc } m\ c\ (\text{bin-sc } n\ b\ w) = \text{bin-sc } n\ b\ (\text{bin-sc } m\ c\ w)$   
 ⟨proof⟩

**lemma** *bin-nth-sc-gen*:

$!!w\ m. \text{bin-nth } (\text{bin-sc } n\ b\ w)\ m = (\text{if } m = n \text{ then } b = \text{bit.B1} \text{ else } \text{bin-nth } w\ m)$   
 $\langle \text{proof} \rangle$

**lemma** *bin-sc-nth [simp]*:

$!!w. (\text{bin-sc } n\ (\text{If } (\text{bin-nth } w\ n)\ \text{bit.B1}\ \text{bit.B0})\ w) = w$   
 $\langle \text{proof} \rangle$

**lemma** *bin-sign-sc [simp]*:

$!!w. \text{bin-sign } (\text{bin-sc } n\ b\ w) = \text{bin-sign } w$   
 $\langle \text{proof} \rangle$

**lemma** *bin-sc-bintr [simp]*:

$!!w\ m. \text{bintrunc } m\ (\text{bin-sc } n\ x\ (\text{bintrunc } m\ (w))) = \text{bintrunc } m\ (\text{bin-sc } n\ x\ w)$   
 $\langle \text{proof} \rangle$

**lemma** *bin-clr-le*:

$!!w. \text{bin-sc } n\ \text{bit.B0}\ w \leq w$   
 $\langle \text{proof} \rangle$

**lemma** *bin-set-ge*:

$!!w. \text{bin-sc } n\ \text{bit.B1}\ w \geq w$   
 $\langle \text{proof} \rangle$

**lemma** *bintr-bin-clr-le*:

$!!w\ m. \text{bintrunc } n\ (\text{bin-sc } m\ \text{bit.B0}\ w) \leq \text{bintrunc } n\ w$   
 $\langle \text{proof} \rangle$

**lemma** *bintr-bin-set-ge*:

$!!w\ m. \text{bintrunc } n\ (\text{bin-sc } m\ \text{bit.B1}\ w) \geq \text{bintrunc } n\ w$   
 $\langle \text{proof} \rangle$

**lemma** *bin-sc-FP [simp]*:  $\text{bin-sc } n\ \text{bit.B0}\ \text{Int.Pl} = \text{Int.Pl}$

$\langle \text{proof} \rangle$

**lemma** *bin-sc-TM [simp]*:  $\text{bin-sc } n\ \text{bit.B1}\ \text{Int.Min} = \text{Int.Min}$

$\langle \text{proof} \rangle$

**lemmas** *bin-sc-simps* = *bin-sc.Z bin-sc.Suc bin-sc-TM bin-sc-FP*

**lemma** *bin-sc-minus*:

$0 < n \implies \text{bin-sc } (\text{Suc } (n - 1))\ b\ w = \text{bin-sc } n\ b\ w$   
 $\langle \text{proof} \rangle$

**lemmas** *bin-sc-Suc-minus* =

*trans [OF bin-sc-minus [symmetric] bin-sc.Suc, standard]*

**lemmas** *bin-sc-Suc-pred [simp]* =

*bin-sc-Suc-minus [of number-of bin, simplified nobm1, standard]*



### 6.3 Operations on lists of booleans

**primrec** *bl-to-bin-aux* :: *bool list*  $\Rightarrow$  *int*  $\Rightarrow$  *int* **where**  
*Nil*: *bl-to-bin-aux* [] *w* = *w*  
| *Cons*: *bl-to-bin-aux* (*b* # *bs*) *w* =  
*bl-to-bin-aux* *bs* (*w* BIT (if *b* then bit.B1 else bit.B0))

**definition** *bl-to-bin* :: *bool list*  $\Rightarrow$  *int* **where**  
*bl-to-bin-def* : *bl-to-bin* *bs* = *bl-to-bin-aux* *bs* *Int.Pls*

**primrec** *bin-to-bl-aux* :: *nat*  $\Rightarrow$  *int*  $\Rightarrow$  *bool list*  $\Rightarrow$  *bool list* **where**  
*Z*: *bin-to-bl-aux* 0 *w* *bl* = *bl*  
| *Suc*: *bin-to-bl-aux* (*Suc* *n*) *w* *bl* =  
*bin-to-bl-aux* *n* (*bin-rest* *w*) ((*bin-last* *w* = bit.B1) # *bl*)

**definition** *bin-to-bl* :: *nat*  $\Rightarrow$  *int*  $\Rightarrow$  *bool list* **where**  
*bin-to-bl-def* : *bin-to-bl* *n* *w* = *bin-to-bl-aux* *n* *w* []

**primrec** *bl-of-nth* :: *nat*  $\Rightarrow$  (*nat*  $\Rightarrow$  *bool*)  $\Rightarrow$  *bool list* **where**  
*Suc*: *bl-of-nth* (*Suc* *n*) *f* = *f* *n* # *bl-of-nth* *n* *f*  
| *Z*: *bl-of-nth* 0 *f* = []

**primrec** *takefill* :: '*a*  $\Rightarrow$  *nat*  $\Rightarrow$  '*a* list  $\Rightarrow$  '*a* list **where**  
*Z*: *takefill* *fill* 0 *xs* = []  
| *Suc*: *takefill* *fill* (*Suc* *n*) *xs* = (  
*case* *xs* of [] => *fill* # *takefill* *fill* *n* *xs*  
| *y* # *ys* => *y* # *takefill* *fill* *n* *ys*)

**definition** *map2* :: ('*a*  $\Rightarrow$  '*b*  $\Rightarrow$  '*c*)  $\Rightarrow$  '*a* list  $\Rightarrow$  '*b* list  $\Rightarrow$  '*c* list **where**  
*map2* *f* *as* *bs* = *map* (*split* *f*) (*zip* *as* *bs*)

### 6.4 Splitting and concatenation

**definition** *bin-rcat* :: *nat*  $\Rightarrow$  *int* list  $\Rightarrow$  *int* **where**  
*bin-rcat* *n* = *foldl* (%*u v. bin-cat* *u* *n* *v*) *Int.Pls*

**function** *bin-rsplit-aux* :: *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  *int*  $\Rightarrow$  *int* list  $\Rightarrow$  *int* list **where**  
*bin-rsplit-aux* *n* *m* *c* *bs* =  
(if *m* = 0 | *n* = 0 then *bs* else  
let (*a*, *b*) = *bin-split* *n* *c*  
in *bin-rsplit-aux* *n* (*m* - *n*) *a* (*b* # *bs*))  
<proof>  
**termination** <proof>

**definition** *bin-rsplit* :: *nat*  $\Rightarrow$  *nat*  $\times$  *int*  $\Rightarrow$  *int* list **where**  
*bin-rsplit* *n* *w* = *bin-rsplit-aux* *n* (*fst* *w*) (*snd* *w*) []

**function** *bin-rsplittl-aux* :: *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  *int*  $\Rightarrow$  *int* list  $\Rightarrow$  *int* list **where**  
*bin-rsplittl-aux* *n* *m* *c* *bs* =  
(if *m* = 0 | *n* = 0 then *bs* else

$\text{let } (a, b) = \text{bin-split } (\text{min } m \ n) \ c$   
 $\text{in bin-rsplittl-aux } n \ (m - n) \ a \ (b \ \# \ bs))$   
 $\langle \text{proof} \rangle$   
**termination**  $\langle \text{proof} \rangle$

**definition**  $\text{bin-rsplittl} :: \text{nat} \Rightarrow \text{nat} \times \text{int} \Rightarrow \text{int list}$  **where**  
 $\text{bin-rsplittl } n \ w = \text{bin-rsplittl-aux } n \ (\text{fst } w) \ (\text{snd } w) \ []$

**declare**  $\text{bin-rsplit-aux.simps}$   $[\text{simp del}]$   
**declare**  $\text{bin-rsplittl-aux.simps}$   $[\text{simp del}]$

**lemma**  $\text{bin-sign-cat}$ :  
 $!!y. \text{bin-sign } (\text{bin-cat } x \ n \ y) = \text{bin-sign } x$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bin-cat-Suc-Bit}$ :  
 $\text{bin-cat } w \ (\text{Suc } n) \ (v \ \text{BIT } b) = \text{bin-cat } w \ n \ v \ \text{BIT } b$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bin-nth-cat}$ :  
 $!!n \ y. \text{bin-nth } (\text{bin-cat } x \ k \ y) \ n =$   
 $(\text{if } n < k \text{ then } \text{bin-nth } y \ n \text{ else } \text{bin-nth } x \ (n - k))$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bin-nth-split}$ :  
 $!!b \ c. \text{bin-split } n \ c = (a, b) ==>$   
 $(\text{ALL } k. \text{bin-nth } a \ k = \text{bin-nth } c \ (n + k)) \ \&$   
 $(\text{ALL } k. \text{bin-nth } b \ k = (k < n \ \& \ \text{bin-nth } c \ k))$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bin-cat-assoc}$ :  
 $!!z. \text{bin-cat } (\text{bin-cat } x \ m \ y) \ n \ z = \text{bin-cat } x \ (m + n) \ (\text{bin-cat } y \ n \ z)$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bin-cat-assoc-sym}$ :  $!!z \ m.$   
 $\text{bin-cat } x \ m \ (\text{bin-cat } y \ n \ z) = \text{bin-cat } (\text{bin-cat } x \ (m - n) \ y) \ (\text{min } m \ n) \ z$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bin-cat-Pls}$   $[\text{simp}]$ :  
 $!!w. \text{bin-cat } \text{Int.Pls } n \ w = \text{bintrunc } n \ w$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bintr-cat1}$ :  
 $!!b. \text{bintrunc } (k + n) \ (\text{bin-cat } a \ n \ b) = \text{bin-cat } (\text{bintrunc } k \ a) \ n \ b$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bintr-cat}$ :  $\text{bintrunc } m \ (\text{bin-cat } a \ n \ b) =$   
 $\text{bin-cat } (\text{bintrunc } (m - n) \ a) \ n \ (\text{bintrunc } (\text{min } m \ n) \ b)$   
 $\langle \text{proof} \rangle$

**lemma** *bintr-cat-same* [simp]:

$$\text{bintrunc } n \text{ (bin-cat } a \text{ } n \text{ } b) = \text{bintrunc } n \text{ } b$$

⟨proof⟩

**lemma** *cat-bintr* [simp]:

$$!!b. \text{bin-cat } a \text{ } n \text{ (bintrunc } n \text{ } b) = \text{bin-cat } a \text{ } n \text{ } b$$

⟨proof⟩

**lemma** *split-bintrunc*:

$$!!b \text{ } c. \text{bin-split } n \text{ } c = (a, b) ==> b = \text{bintrunc } n \text{ } c$$

⟨proof⟩

**lemma** *bin-cat-split*:

$$!!v \text{ } w. \text{bin-split } n \text{ } w = (u, v) ==> w = \text{bin-cat } u \text{ } n \text{ } v$$

⟨proof⟩

**lemma** *bin-split-cat*:

$$!!w. \text{bin-split } n \text{ (bin-cat } v \text{ } n \text{ } w) = (v, \text{bintrunc } n \text{ } w)$$

⟨proof⟩

**lemma** *bin-split-Pls* [simp]:

$$\text{bin-split } n \text{ Int.Pls} = (\text{Int.Pls}, \text{Int.Pls})$$

⟨proof⟩

**lemma** *bin-split-Min* [simp]:

$$\text{bin-split } n \text{ Int.Min} = (\text{Int.Min}, \text{bintrunc } n \text{ Int.Min})$$

⟨proof⟩

**lemma** *bin-split-trunc*:

$$!!m \text{ } b \text{ } c. \text{bin-split (min } m \text{ } n) \text{ } c = (a, b) ==>$$

$$\text{bin-split } n \text{ (bintrunc } m \text{ } c) = (\text{bintrunc (} m - n \text{)} \text{ } a, b)$$

⟨proof⟩

**lemma** *bin-split-trunc1*:

$$!!m \text{ } b \text{ } c. \text{bin-split } n \text{ } c = (a, b) ==>$$

$$\text{bin-split } n \text{ (bintrunc } m \text{ } c) = (\text{bintrunc (} m - n \text{)} \text{ } a, \text{bintrunc } m \text{ } b)$$

⟨proof⟩

**lemma** *bin-cat-num*:

$$!!b. \text{bin-cat } a \text{ } n \text{ } b = a * 2 ^ n + \text{bintrunc } n \text{ } b$$

⟨proof⟩

**lemma** *bin-split-num*:

$$!!b. \text{bin-split } n \text{ } b = (b \text{ div } 2 ^ n, b \text{ mod } 2 ^ n)$$

⟨proof⟩

## 6.5 Miscellaneous lemmas

**lemma** *nth-2p-bin*:

!!*m. bin-nth* ( $2 \wedge n$ ) *m* = (*m* = *n*)  
 ⟨*proof*⟩

**lemma** *ex-eq-or*:

(*EX m. n* = *Suc m* & (*m* = *k* | *P m*)) = (*n* = *Suc k* | (*EX m. n* = *Suc m* & *P m*))  
 ⟨*proof*⟩

**end**

## 7 BinBoolList: Bool lists and integers

**theory** *BinBoolList*

**imports** *BinOperations*

**begin**

### 7.1 Arithmetic in terms of bool lists

**primrec** *rbl-succ* :: *bool list* => *bool list* **where**

*Nil*: *rbl-succ Nil* = *Nil*  
 | *Cons*: *rbl-succ* (*x* # *xs*) = (if *x* then *False* # *rbl-succ xs* else *True* # *xs*)

**primrec** *rbl-pred* :: *bool list* => *bool list* **where**

*Nil*: *rbl-pred Nil* = *Nil*  
 | *Cons*: *rbl-pred* (*x* # *xs*) = (if *x* then *False* # *xs* else *True* # *rbl-pred xs*)

**primrec** *rbl-add* :: *bool list* => *bool list* => *bool list* **where**

*Nil*: *rbl-add Nil x* = *Nil*  
 | *Cons*: *rbl-add* (*y* # *ys*) *x* = (let *ws* = *rbl-add ys* (*tl x*) in  
   (*y* ~ = *hd x*) # (if *hd x* & *y* then *rbl-succ ws* else *ws*))

**primrec** *rbl-mult* :: *bool list* => *bool list* => *bool list* **where**

*Nil*: *rbl-mult Nil x* = *Nil*  
 | *Cons*: *rbl-mult* (*y* # *ys*) *x* = (let *ws* = *False* # *rbl-mult ys x* in  
   if *y* then *rbl-add ws x* else *ws*)

**lemma** *butlast-power*:

(*butlast*  $\wedge n$ ) *bl* = *take* (*length bl* - *n*) *bl*  
 ⟨*proof*⟩

**lemma** *bin-to-bl-aux-Pls-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \text{ Int.Pls } bl = \\ \text{bin-to-bl-aux } (n - 1) \text{ Int.Pls } (\text{False} \# bl) \\ \langle \text{proof} \rangle$$

**lemma** *bin-to-bl-aux-Min-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \text{ Int.Min } bl = \\ \text{bin-to-bl-aux } (n - 1) \text{ Int.Min } (\text{True} \# bl) \\ \langle \text{proof} \rangle$$

**lemma** *bin-to-bl-aux-Bit-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \text{ (w BIT b) } bl = \\ \text{bin-to-bl-aux } (n - 1) \text{ w } ((b = \text{bit.B1}) \# bl) \\ \langle \text{proof} \rangle$$

**lemma** *bin-to-bl-aux-Bit0-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \text{ (Int.Bit0 w) } bl = \\ \text{bin-to-bl-aux } (n - 1) \text{ w } (\text{False} \# bl) \\ \langle \text{proof} \rangle$$

**lemma** *bin-to-bl-aux-Bit1-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \text{ (Int.Bit1 w) } bl = \\ \text{bin-to-bl-aux } (n - 1) \text{ w } (\text{True} \# bl) \\ \langle \text{proof} \rangle$$

**lemma** *bl-to-bin-aux-append*:

$$\text{bl-to-bin-aux } (bs @ cs) \text{ w} = \text{bl-to-bin-aux } cs \text{ (bl-to-bin-aux } bs \text{ w)} \\ \langle \text{proof} \rangle$$

**lemma** *bin-to-bl-aux-append*:

$$\text{bin-to-bl-aux } n \text{ w } bs @ cs = \text{bin-to-bl-aux } n \text{ w } (bs @ cs) \\ \langle \text{proof} \rangle$$

**lemma** *bl-to-bin-append*:

$$\text{bl-to-bin } (bs @ cs) = \text{bl-to-bin-aux } cs \text{ (bl-to-bin } bs) \\ \langle \text{proof} \rangle$$

**lemma** *bin-to-bl-aux-alt*:

$$\text{bin-to-bl-aux } n \text{ w } bs = \text{bin-to-bl } n \text{ w } @ bs \\ \langle \text{proof} \rangle$$

**lemma** *bin-to-bl-0*:  $\text{bin-to-bl } 0 \text{ } bs = []$

$\langle \text{proof} \rangle$

**lemma** *size-bin-to-bl-aux*:

$$\text{size } (\text{bin-to-bl-aux } n \text{ w } bs) = n + \text{length } bs \\ \langle \text{proof} \rangle$$

**lemma** *size-bin-to-bl*:  $\text{size } (\text{bin-to-bl } n \ w) = n$   
 ⟨proof⟩

**lemma** *bin-bl-bin'*:  
 $\text{bl-to-bin } (\text{bin-to-bl-aux } n \ w \ bs) =$   
 $\text{bl-to-bin-aux } bs \ (\text{bintrunc } n \ w)$   
 ⟨proof⟩

**lemma** *bin-bl-bin*:  $\text{bl-to-bin } (\text{bin-to-bl } n \ w) = \text{bintrunc } n \ w$   
 ⟨proof⟩

**lemma** *bl-bin-bl'*:  
 $\text{bin-to-bl } (n + \text{length } bs) \ (\text{bl-to-bin-aux } bs \ w) =$   
 $\text{bin-to-bl-aux } n \ w \ bs$   
 ⟨proof⟩

**lemma** *bl-bin-bl*:  $\text{bin-to-bl } (\text{length } bs) \ (\text{bl-to-bin } bs) = bs$   
 ⟨proof⟩

**declare**  
 $\text{bin-to-bl-0} \ [\text{simp}]$   
 $\text{size-bin-to-bl} \ [\text{simp}]$   
 $\text{bin-bl-bin} \ [\text{simp}]$   
 $\text{bl-bin-bl} \ [\text{simp}]$

**lemma** *bl-to-bin-inj*:  
 $\text{bl-to-bin } bs = \text{bl-to-bin } cs \implies \text{length } bs = \text{length } cs \implies bs = cs$   
 ⟨proof⟩

**lemma** *bl-to-bin-False*:  $\text{bl-to-bin } (\text{False} \ \# \ bl) = \text{bl-to-bin } bl$   
 ⟨proof⟩

**lemma** *bl-to-bin-Nil*:  $\text{bl-to-bin } [] = \text{Int.Pls}$   
 ⟨proof⟩

**lemma** *bin-to-bl-Pls-aux*:  
 $\text{bin-to-bl-aux } n \ \text{Int.Pls } bl = \text{replicate } n \ \text{False} \ @ \ bl$   
 ⟨proof⟩

**lemma** *bin-to-bl-Pls*:  $\text{bin-to-bl } n \ \text{Int.Pls} = \text{replicate } n \ \text{False}$   
 ⟨proof⟩

**lemma** *bin-to-bl-Min-aux* [rule-format] :  
 $\text{ALL } bl. \ \text{bin-to-bl-aux } n \ \text{Int.Min } bl = \text{replicate } n \ \text{True} \ @ \ bl$   
 ⟨proof⟩

**lemma** *bin-to-bl-Min*:  $\text{bin-to-bl } n \ \text{Int.Min} = \text{replicate } n \ \text{True}$   
 ⟨proof⟩

**lemma** *bl-to-bin-rep-F*:

*bl-to-bin (replicate n False @ bl) = bl-to-bin bl*  
*<proof>*

**lemma** *bin-to-bl-trunc*:

*n <= m ==> bin-to-bl n (bintrunc m w) = bin-to-bl n w*  
*<proof>*

**declare**

*bin-to-bl-trunc [simp]*  
*bl-to-bin-False [simp]*  
*bl-to-bin-Nil [simp]*

**lemma** *bin-to-bl-aux-bintr [rule-format]* :

*ALL m bin bl. bin-to-bl-aux n (bintrunc m bin) bl =*  
*replicate (n - m) False @ bin-to-bl-aux (min n m) bin bl*  
*<proof>*

**lemmas** *bin-to-bl-bintr =*

*bin-to-bl-aux-bintr [where bl = [], folded bin-to-bl-def]*

**lemma** *bl-to-bin-rep-False*: *bl-to-bin (replicate n False) = Int.Pls*

*<proof>*

**lemma** *len-bin-to-bl-aux*:

*length (bin-to-bl-aux n w bs) = n + length bs*  
*<proof>*

**lemma** *len-bin-to-bl [simp]*: *length (bin-to-bl n w) = n*

*<proof>*

**lemma** *sign-bl-bin'*:

*bin-sign (bl-to-bin-aux bs w) = bin-sign w*  
*<proof>*

**lemma** *sign-bl-bin*: *bin-sign (bl-to-bin bs) = Int.Pls*

*<proof>*

**lemma** *bl-sbin-sign-aux*:

*hd (bin-to-bl-aux (Suc n) w bs) =*  
*(bin-sign (sbintrunc n w) = Int.Min)*  
*<proof>*

**lemma** *bl-sbin-sign*:

*hd (bin-to-bl (Suc n) w) = (bin-sign (sbintrunc n w) = Int.Min)*  
*<proof>*

**lemma** *bin-nth-of-bl-aux [rule-format]*:

$\forall w. \text{bin-nth } (\text{bl-to-bin-aux } \text{bl } w) \text{ } n =$   
 $(n < \text{size } \text{bl} \ \& \ \text{rev } \text{bl} ! n \mid n \geq \text{length } \text{bl} \ \& \ \text{bin-nth } w \ (n - \text{size } \text{bl}))$   
 $\langle \text{proof} \rangle$

**lemma** *bin-nth-of-bl*:  $\text{bin-nth } (\text{bl-to-bin } \text{bl}) \text{ } n = (n < \text{length } \text{bl} \ \& \ \text{rev } \text{bl} ! n)$   
 $\langle \text{proof} \rangle$

**lemma** *bin-nth-bl* [rule-format] :  $\text{ALL } m \ w. \ n < m \ \longrightarrow$   
 $\text{bin-nth } w \ n = \text{nth } (\text{rev } (\text{bin-to-bl } m \ w)) \ n$   
 $\langle \text{proof} \rangle$

**lemma** *nth-rev* [rule-format] :  
 $n < \text{length } \text{xs} \ \longrightarrow \ \text{rev } \text{xs} ! n = \text{xs} ! (\text{length } \text{xs} - 1 - n)$   
 $\langle \text{proof} \rangle$

**lemmas** *nth-rev-alt* = *nth-rev* [where  $\text{xs} = \text{rev } \text{ys}$ , simplified, standard]

**lemma** *nth-bin-to-bl-aux* [rule-format] :  
 $\text{ALL } w \ n \ \text{bl}. \ n < m + \text{length } \text{bl} \ \longrightarrow \ (\text{bin-to-bl-aux } m \ w \ \text{bl}) ! n =$   
 $(\text{if } n < m \ \text{then } \text{bin-nth } w \ (m - 1 - n) \ \text{else } \text{bl} ! (n - m))$   
 $\langle \text{proof} \rangle$

**lemma** *nth-bin-to-bl*:  $n < m \ \Longrightarrow \ (\text{bin-to-bl } m \ w) ! n = \text{bin-nth } w \ (m - \text{Suc } n)$   
 $\langle \text{proof} \rangle$

**lemma** *bl-to-bin-lt2p-aux* [rule-format]:  
 $\forall w. \ \text{bl-to-bin-aux } \text{bs } w < (w + 1) * (2 ^ \text{length } \text{bs})$   
 $\langle \text{proof} \rangle$

**lemma** *bl-to-bin-lt2p*:  $\text{bl-to-bin } \text{bs} < (2 ^ \text{length } \text{bs})$   
 $\langle \text{proof} \rangle$

**lemma** *bl-to-bin-ge2p-aux* [rule-format] :  
 $\forall w. \ \text{bl-to-bin-aux } \text{bs } w \geq w * (2 ^ \text{length } \text{bs})$   
 $\langle \text{proof} \rangle$

**lemma** *bl-to-bin-ge0*:  $\text{bl-to-bin } \text{bs} \geq 0$   
 $\langle \text{proof} \rangle$

**lemma** *butlast-rest-bin*:  
 $\text{butlast } (\text{bin-to-bl } n \ w) = \text{bin-to-bl } (n - 1) \ (\text{bin-rest } w)$   
 $\langle \text{proof} \rangle$

**lemmas** *butlast-bin-rest* = *butlast-rest-bin*  
[where  $w = \text{bl-to-bin } \text{bl}$  and  $n = \text{length } \text{bl}$ , simplified, standard]

**lemma** *butlast-rest-bl2bin-aux*:  
 $\text{bl} \sim = [] \ \Longrightarrow$   
 $\text{bl-to-bin-aux } (\text{butlast } \text{bl}) \ w = \text{bin-rest } (\text{bl-to-bin-aux } \text{bl } w)$



$\langle \text{proof} \rangle$

**lemma** *butlast-rest-bl2bin*:

$\text{bl-to-bin } (\text{butlast } \text{bl}) = \text{bin-rest } (\text{bl-to-bin } \text{bl})$

$\langle \text{proof} \rangle$

**lemma** *trunc-bl2bin-aux* [rule-format]:

$\text{ALL } w. \text{bintrunc } m \text{ (bl-to-bin-aux } \text{bl } w) =$

$\text{bl-to-bin-aux } (\text{drop } (\text{length } \text{bl} - m) \text{ bl}) (\text{bintrunc } (m - \text{length } \text{bl}) w)$

$\langle \text{proof} \rangle$

**lemma** *trunc-bl2bin*:

$\text{bintrunc } m \text{ (bl-to-bin } \text{bl}) = \text{bl-to-bin } (\text{drop } (\text{length } \text{bl} - m) \text{ bl})$

$\langle \text{proof} \rangle$

**lemmas** *trunc-bl2bin-len* [simp] =

*trunc-bl2bin* [of length bl bl, simplified, standard]

**lemma** *bl2bin-drop*:

$\text{bl-to-bin } (\text{drop } k \text{ bl}) = \text{bintrunc } (\text{length } \text{bl} - k) \text{ (bl-to-bin } \text{bl})$

$\langle \text{proof} \rangle$

**lemma** *nth-rest-power-bin* [rule-format] :

$\text{ALL } n. \text{bin-nth } ((\text{bin-rest } ^k) w) n = \text{bin-nth } w (n + k)$

$\langle \text{proof} \rangle$

**lemma** *take-rest-power-bin*:

$m \leq n \implies \text{take } m \text{ (bin-to-bl } n w) = \text{bin-to-bl } m \text{ ((bin-rest } ^{(n-m)}) w)$

$\langle \text{proof} \rangle$

**lemma** *hd-butlast*:  $\text{size } xs > 1 \implies \text{hd } (\text{butlast } xs) = \text{hd } xs$

$\langle \text{proof} \rangle$

**lemma** *last-bin-last'*:

$\text{size } xs > 0 \implies \text{last } xs = (\text{bin-last } (\text{bl-to-bin-aux } xs w) = \text{bit.B1})$

$\langle \text{proof} \rangle$

**lemma** *last-bin-last*:

$\text{size } xs > 0 \implies \text{last } xs = (\text{bin-last } (\text{bl-to-bin } xs) = \text{bit.B1})$

$\langle \text{proof} \rangle$

**lemma** *bin-last-last*:

$\text{bin-last } w = (\text{if last } (\text{bin-to-bl } (\text{Suc } n) w) \text{ then bit.B1 else bit.B0})$

$\langle \text{proof} \rangle$

**lemma** *map2-Nil* [simp]:  $\text{map2 } f [] ys = []$

$\langle \text{proof} \rangle$

**lemma** *map2-Cons* [*simp*]:  
 $\text{map2 } f \ (x \# \text{xs}) \ (y \# \text{ys}) = f \ x \ y \# \text{map2 } f \ \text{xs} \ \text{ys}$   
 ⟨*proof*⟩

**lemma** *bl-xor-aux-bin* [*rule-format*] : *ALL*  $v \ w \ bs \ cs$ .  
 $\text{map2 } (\%x \ y. \ x \ \sim = \ y) \ (\text{bin-to-bl-aux } n \ v \ bs) \ (\text{bin-to-bl-aux } n \ w \ cs) =$   
 $\text{bin-to-bl-aux } n \ (v \ XOR \ w) \ (\text{map2 } (\%x \ y. \ x \ \sim = \ y) \ bs \ cs)$   
 ⟨*proof*⟩

**lemma** *bl-or-aux-bin* [*rule-format*] : *ALL*  $v \ w \ bs \ cs$ .  
 $\text{map2 } (op \ | \ ) \ (\text{bin-to-bl-aux } n \ v \ bs) \ (\text{bin-to-bl-aux } n \ w \ cs) =$   
 $\text{bin-to-bl-aux } n \ (v \ OR \ w) \ (\text{map2 } (op \ | \ ) \ bs \ cs)$   
 ⟨*proof*⟩

**lemma** *bl-and-aux-bin* [*rule-format*] : *ALL*  $v \ w \ bs \ cs$ .  
 $\text{map2 } (op \ \& \ ) \ (\text{bin-to-bl-aux } n \ v \ bs) \ (\text{bin-to-bl-aux } n \ w \ cs) =$   
 $\text{bin-to-bl-aux } n \ (v \ AND \ w) \ (\text{map2 } (op \ \& \ ) \ bs \ cs)$   
 ⟨*proof*⟩

**lemma** *bl-not-aux-bin* [*rule-format*] :  
*ALL*  $w \ cs$ .  $\text{map } Not \ (\text{bin-to-bl-aux } n \ w \ cs) =$   
 $\text{bin-to-bl-aux } n \ (NOT \ w) \ (\text{map } Not \ cs)$   
 ⟨*proof*⟩

**lemmas** *bl-not-bin* = *bl-not-aux-bin*  
 [where  $cs = []$ , *unfolded bin-to-bl-def* [*symmetric*] *map.simps*]

**lemmas** *bl-and-bin* = *bl-and-aux-bin* [where  $bs=[]$  and  $cs=[]$ ,  
*unfolded map2-Nil, folded bin-to-bl-def*]

**lemmas** *bl-or-bin* = *bl-or-aux-bin* [where  $bs=[]$  and  $cs=[]$ ,  
*unfolded map2-Nil, folded bin-to-bl-def*]

**lemmas** *bl-xor-bin* = *bl-xor-aux-bin* [where  $bs=[]$  and  $cs=[]$ ,  
*unfolded map2-Nil, folded bin-to-bl-def*]

**lemma** *drop-bin2bl-aux* [*rule-format*] :  
*ALL*  $m \ bin \ bs$ .  $\text{drop } m \ (\text{bin-to-bl-aux } n \ bin \ bs) =$   
 $\text{bin-to-bl-aux } (n - m) \ bin \ (\text{drop } (m - n) \ bs)$   
 ⟨*proof*⟩

**lemma** *drop-bin2bl*:  $\text{drop } m \ (\text{bin-to-bl } n \ bin) = \text{bin-to-bl } (n - m) \ bin$   
 ⟨*proof*⟩

**lemma** *take-bin2bl-lem1* [*rule-format*] :  
*ALL*  $w \ bs$ .  $\text{take } m \ (\text{bin-to-bl-aux } m \ w \ bs) = \text{bin-to-bl } m \ w$   
 ⟨*proof*⟩

**lemma** *take-bin2bl-lem* [rule-format] :

*ALL w bs. take m (bin-to-bl-aux (m + n) w bs) =  
take m (bin-to-bl (m + n) w)  
<proof>*

**lemma** *bin-split-take* [rule-format] :

*ALL b c. bin-split n c = (a, b) -->  
bin-to-bl m a = take m (bin-to-bl (m + n) c)  
<proof>*

**lemma** *bin-split-take1*:

*k = m + n ==> bin-split n c = (a, b) ==>  
bin-to-bl m a = take m (bin-to-bl k c)  
<proof>*

**lemma** *nth-takefill* [rule-format] : *ALL m l. m < n -->*

*takefill fill n l ! m = (if m < length l then l ! m else fill)  
<proof>*

**lemma** *takefill-alt* [rule-format] :

*ALL l. takefill fill n l = take n l @ replicate (n - length l) fill  
<proof>*

**lemma** *takefill-replicate* [simp]:

*takefill fill n (replicate m fill) = replicate n fill  
<proof>*

**lemma** *takefill-le'* [rule-format] :

*ALL l n. n = m + k --> takefill x m (takefill x n l) = takefill x m l  
<proof>*

**lemma** *length-takefill* [simp]: *length (takefill fill n l) = n*

*<proof>*

**lemma** *take-takefill'*:

*!!w n. n = k + m ==> take k (takefill fill n w) = takefill fill k w  
<proof>*

**lemma** *drop-takefill*:

*!!w. drop k (takefill fill (m + k) w) = takefill fill m (drop k w)  
<proof>*

**lemma** *takefill-le* [simp]:

*m ≤ n ==> takefill x m (takefill x n l) = takefill x m l  
<proof>*

**lemma** *take-takefill* [simp]:

*m ≤ n ==> take m (takefill fill n w) = takefill fill m w  
<proof>*

**lemma** *takefill-append*:

$\text{takefill fill } (m + \text{length } xs) \ (xs \ @ \ w) = xs \ @ \ (\text{takefill fill } m \ w)$   
 $\langle \text{proof} \rangle$

**lemma** *takefill-same'*:

$l = \text{length } xs \implies \text{takefill fill } l \ xs = xs$   
 $\langle \text{proof} \rangle$

**lemmas** *takefill-same* [simp] = *takefill-same'* [OF refl]

**lemma** *takefill-bintrunc*:

$\text{takefill False } n \ bl = \text{rev } (\text{bin-to-bl } n \ (\text{bl-to-bin } (\text{rev } bl)))$   
 $\langle \text{proof} \rangle$

**lemma** *bl-bin-bl-rtf*:

$\text{bin-to-bl } n \ (\text{bl-to-bin } bl) = \text{rev } (\text{takefill False } n \ (\text{rev } bl))$   
 $\langle \text{proof} \rangle$

**lemmas** *bl-bin-bl-rep-drop* =

*bl-bin-bl-rtf* [simplified takefill-alt,  
simplified, simplified rev-take, simplified]

**lemma** *tf-rev*:

$n + k = m + \text{length } bl \implies \text{takefill } x \ m \ (\text{rev } (\text{takefill } y \ n \ bl)) =$   
 $\text{rev } (\text{takefill } y \ m \ (\text{rev } (\text{takefill } x \ k \ (\text{rev } bl))))$   
 $\langle \text{proof} \rangle$

**lemma** *takefill-minus*:

$0 < n \implies \text{takefill fill } (\text{Suc } (n - 1)) \ w = \text{takefill fill } n \ w$   
 $\langle \text{proof} \rangle$

**lemmas** *takefill-Suc-cases* =

*list.cases* [THEN *takefill.Suc* [THEN *trans*], *standard*]

**lemmas** *takefill-Suc-Nil* = *takefill-Suc-cases* (1)

**lemmas** *takefill-Suc-Cons* = *takefill-Suc-cases* (2)

**lemmas** *takefill-minus-simps* = *takefill-Suc-cases* [THEN [2]

*takefill-minus* [symmetric, THEN *trans*], *standard*]

**lemmas** *takefill-pred-simps* [simp] =

*takefill-minus-simps* [where *n=number-of bin*, *simplified nobm1*, *standard*]

**lemma** *bl-to-bin-aux-cat*:

$!!nv \ v. \ \text{bl-to-bin-aux } bs \ (\text{bin-cat } w \ nv \ v) =$   
 $\text{bin-cat } w \ (nv + \text{length } bs) \ (\text{bl-to-bin-aux } bs \ v)$

$\langle \text{proof} \rangle$

**lemma** *bin-to-bl-aux-cat*:

$!!w \text{ bs. } \text{bin-to-bl-aux } (nv + nw) (\text{bin-cat } v \text{ nw } w) \text{ bs} =$   
 $\text{bin-to-bl-aux } nv \text{ v } (\text{bin-to-bl-aux } nw \text{ w } \text{bs})$   
 $\langle \text{proof} \rangle$

**lemmas** *bl-to-bin-aux-alt* =

*bl-to-bin-aux-cat* [**where**  $nv = 0$  **and**  $v = \text{Int.Pls}$ ,  
*simplified bl-to-bin-def* [*symmetric*], *simplified*]

**lemmas** *bin-to-bl-cat* =

*bin-to-bl-aux-cat* [**where**  $\text{bs} = []$ , *folded bin-to-bl-def*]

**lemmas** *bl-to-bin-aux-app-cat* =

*trans* [*OF bl-to-bin-aux-append bl-to-bin-aux-alt*]

**lemmas** *bin-to-bl-aux-cat-app* =

*trans* [*OF bin-to-bl-aux-cat bin-to-bl-aux-alt*]

**lemmas** *bl-to-bin-app-cat* = *bl-to-bin-aux-app-cat*

[**where**  $w = \text{Int.Pls}$ , *folded bl-to-bin-def*]

**lemmas** *bin-to-bl-cat-app* = *bin-to-bl-aux-cat-app*

[**where**  $\text{bs} = []$ , *folded bin-to-bl-def*]

**lemma** *bl-to-bin-app-cat-alt*:

$\text{bin-cat } (\text{bl-to-bin } cs) \text{ n } w = \text{bl-to-bin } (cs @ \text{bin-to-bl } n \text{ w})$   
 $\langle \text{proof} \rangle$

**lemma** *mask-lem*:  $(\text{bl-to-bin } (\text{True} \# \text{replicate } n \text{ False})) =$

$\text{Int.succ } (\text{bl-to-bin } (\text{replicate } n \text{ True}))$   
 $\langle \text{proof} \rangle$

**lemma** *length-bl-of-nth* [*simp*]:  $\text{length } (\text{bl-of-nth } n \text{ f}) = n$

$\langle \text{proof} \rangle$

**lemma** *nth-bl-of-nth* [*simp*]:

$m < n \implies \text{rev } (\text{bl-of-nth } n \text{ f}) ! m = f \text{ m}$   
 $\langle \text{proof} \rangle$

**lemma** *bl-of-nth-inj*:

$(!!k. k < n \implies f \text{ k} = g \text{ k}) \implies \text{bl-of-nth } n \text{ f} = \text{bl-of-nth } n \text{ g}$   
 $\langle \text{proof} \rangle$

**lemma** *bl-of-nth-nth-le* [*rule-format*] : *ALL xs.*

$\text{length } xs \geq n \implies \text{bl-of-nth } n (\text{nth } (\text{rev } xs)) = \text{drop } (\text{length } xs - n) \text{ xs}$

$\langle proof \rangle$

**lemmas** *bl-of-nth-nth* [*simp*] = *order-refl* [*THEN bl-of-nth-nth-le, simplified*]

**lemma** *size-rbl-pred*: *length (rbl-pred bl) = length bl*  
 $\langle proof \rangle$

**lemma** *size-rbl-succ*: *length (rbl-succ bl) = length bl*  
 $\langle proof \rangle$

**lemma** *size-rbl-add*:  
 $!!cl. \text{length (rbl-add bl cl) = length bl}$   
 $\langle proof \rangle$

**lemma** *size-rbl-mult*:  
 $!!cl. \text{length (rbl-mult bl cl) = length bl}$   
 $\langle proof \rangle$

**lemmas** *rbl-sizes* [*simp*] =  
*size-rbl-pred size-rbl-succ size-rbl-add size-rbl-mult*

**lemmas** *rbl-Nils* =  
*rbl-pred.Nil rbl-succ.Nil rbl-add.Nil rbl-mult.Nil*

**lemma** *rbl-pred*:  
 $!!bin. \text{rbl-pred (rev (bin-to-bl n bin)) = rev (bin-to-bl n (Int.pred bin))}$   
 $\langle proof \rangle$

**lemma** *rbl-succ*:  
 $!!bin. \text{rbl-succ (rev (bin-to-bl n bin)) = rev (bin-to-bl n (Int.succ bin))}$   
 $\langle proof \rangle$

**lemma** *rbl-add*:  
 $!!bina \ binb. \text{rbl-add (rev (bin-to-bl n bina)) (rev (bin-to-bl n binb)) =}$   
 $\text{rev (bin-to-bl n (bina + binb))}$   
 $\langle proof \rangle$

**lemma** *rbl-add-app2*:  
 $!!blb. \text{length blb} \geq \text{length bla} \implies$   
 $\text{rbl-add bla (blb @ blc) = rbl-add bla blb}$   
 $\langle proof \rangle$

**lemma** *rbl-add-take2*:  
 $!!blb. \text{length blb} \geq \text{length bla} \implies$   
 $\text{rbl-add bla (take (length bla) blb) = rbl-add bla blb}$   
 $\langle proof \rangle$

**lemma** *rbl-add-long*:  
 $m \geq n \implies \text{rbl-add (rev (bin-to-bl n bina)) (rev (bin-to-bl m binb)) =}$

$rev (bin-to-bl\ n\ (bina + binb))$   
 $\langle proof \rangle$

**lemma** *rbl-mult-app2*:  
 $!!blb. length\ blb \geq length\ bla \implies$   
 $rbl-mult\ bla\ (blb @ blc) = rbl-mult\ bla\ blb$   
 $\langle proof \rangle$

**lemma** *rbl-mult-take2*:  
 $length\ blb \geq length\ bla \implies$   
 $rbl-mult\ bla\ (take\ (length\ bla)\ blb) = rbl-mult\ bla\ blb$   
 $\langle proof \rangle$

**lemma** *rbl-mult-gt1*:  
 $m \geq length\ bl \implies rbl-mult\ bl\ (rev\ (bin-to-bl\ m\ binb)) =$   
 $rbl-mult\ bl\ (rev\ (bin-to-bl\ (length\ bl)\ binb))$   
 $\langle proof \rangle$

**lemma** *rbl-mult-gt*:  
 $m > n \implies rbl-mult\ (rev\ (bin-to-bl\ n\ bina))\ (rev\ (bin-to-bl\ m\ binb)) =$   
 $rbl-mult\ (rev\ (bin-to-bl\ n\ bina))\ (rev\ (bin-to-bl\ n\ binb))$   
 $\langle proof \rangle$

**lemmas** *rbl-mult-Suc* = *lessI* [THEN *rbl-mult-gt*]

**lemma** *rbbl-Cons*:  
 $b \# rev\ (bin-to-bl\ n\ x) = rev\ (bin-to-bl\ (Suc\ n)\ (x\ BIT\ If\ b\ bit.B1\ bit.B0))$   
 $\langle proof \rangle$

**lemma** *rbl-mult*:  $!!bina\ binb.$   
 $rbl-mult\ (rev\ (bin-to-bl\ n\ bina))\ (rev\ (bin-to-bl\ n\ binb)) =$   
 $rev\ (bin-to-bl\ n\ (bina * binb))$   
 $\langle proof \rangle$

**lemma** *rbl-add-split*:  
 $P\ (rbl-add\ (y \# ys)\ (x \# xs)) =$   
 $(ALL\ ws. length\ ws = length\ ys \longrightarrow ws = rbl-add\ ys\ xs \longrightarrow$   
 $(y \longrightarrow ((x \longrightarrow P\ (False \# rbl-succ\ ws)) \ \&\ (\sim x \longrightarrow P\ (True \# ws)))) \ \&$   
 $(\sim y \longrightarrow P\ (x \# ws)))$   
 $\langle proof \rangle$

**lemma** *rbl-mult-split*:  
 $P\ (rbl-mult\ (y \# ys)\ xs) =$   
 $(ALL\ ws. length\ ws = Suc\ (length\ ys) \longrightarrow ws = False \# rbl-mult\ ys\ xs \longrightarrow$   
 $(y \longrightarrow P\ (rbl-add\ ws\ xs)) \ \&\ (\sim y \longrightarrow P\ ws))$   
 $\langle proof \rangle$

**lemma** *and-len*:  $xs = ys \implies xs = ys \ \&\ length\ xs = length\ ys$

$\langle \text{proof} \rangle$

**lemma** *size-if*:  $\text{size } (\text{if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then } \text{size } xs \text{ else } \text{size } ys)$   
 $\langle \text{proof} \rangle$

**lemma** *tl-if*:  $\text{tl } (\text{if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then } \text{tl } xs \text{ else } \text{tl } ys)$   
 $\langle \text{proof} \rangle$

**lemma** *hd-if*:  $\text{hd } (\text{if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then } \text{hd } xs \text{ else } \text{hd } ys)$   
 $\langle \text{proof} \rangle$

**lemma** *if-Not-x*:  $(\text{if } p \text{ then } \sim x \text{ else } x) = (p = (\sim x))$   
 $\langle \text{proof} \rangle$

**lemma** *if-x-Not*:  $(\text{if } p \text{ then } x \text{ else } \sim x) = (p = x)$   
 $\langle \text{proof} \rangle$

**lemma** *if-same-and*:  $(\text{If } p \text{ } x \text{ } y \ \& \ \text{If } p \text{ } u \text{ } v) = (\text{if } p \text{ then } x \ \& \ u \text{ else } y \ \& \ v)$   
 $\langle \text{proof} \rangle$

**lemma** *if-same-eq*:  $(\text{If } p \text{ } x \text{ } y \ = \ (\text{If } p \text{ } u \text{ } v)) = (\text{if } p \text{ then } x = (u) \text{ else } y = (v))$   
 $\langle \text{proof} \rangle$

**lemma** *if-same-eq-not*:  
 $(\text{If } p \text{ } x \text{ } y \ = \ (\sim \text{If } p \text{ } u \text{ } v)) = (\text{if } p \text{ then } x = (\sim u) \text{ else } y = (\sim v))$   
 $\langle \text{proof} \rangle$

**lemma** *if-Cons*:  $(\text{if } p \text{ then } x \ \# \ xs \text{ else } y \ \# \ ys) = \text{If } p \text{ } x \text{ } y \ \# \ \text{If } p \text{ } xs \text{ } ys$   
 $\langle \text{proof} \rangle$

**lemma** *if-single*:  
 $(\text{if } xc \text{ then } [xab] \text{ else } [an]) = [\text{if } xc \text{ then } xab \text{ else } an]$   
 $\langle \text{proof} \rangle$

**lemma** *if-bool-simps*:  
 $\text{If } p \text{ } \text{True} \text{ } y = (p \mid y) \ \& \ \text{If } p \text{ } \text{False} \text{ } y = (\sim p \ \& \ y) \ \& \$   
 $\text{If } p \text{ } y \text{ } \text{True} = (p \dashrightarrow y) \ \& \ \text{If } p \text{ } y \text{ } \text{False} = (p \ \& \ y)$   
 $\langle \text{proof} \rangle$

**lemmas** *if-simps* = *if-x-Not if-Not-x if-cancel if-True if-False if-bool-simps*

**lemmas** *seqr* = *eq-reflection* [where  $x = \text{size } w$ , standard]

**lemmas** *tl-Nil* = *tl.simps* (1)

**lemmas** *tl-Cons* = *tl.simps* (2)



## 7.2 Repeated splitting or concatenation

**lemma** *sclem*:

$\text{size } (\text{concat } (\text{map } (\text{bin-to-bl } n) \text{ xs})) = \text{length } \text{xs} * n$   
 $\langle \text{proof} \rangle$

**lemma** *bin-cat-foldl-lem* [rule-format] :

$\text{ALL } x. \text{foldl } (\%u. \text{bin-cat } u \ n) \ x \ \text{xs} =$   
 $\text{bin-cat } x \ (\text{size } \text{xs} * n) \ (\text{foldl } (\%u. \text{bin-cat } u \ n) \ y \ \text{xs})$   
 $\langle \text{proof} \rangle$

**lemma** *bin-rcat-bl*:

$(\text{bin-rcat } n \ \text{wl}) = \text{bl-to-bin } (\text{concat } (\text{map } (\text{bin-to-bl } n) \ \text{wl}))$   
 $\langle \text{proof} \rangle$

**lemmas** *bin-rsplit-aux-simps* = *bin-rsplit-aux.simps* *bin-rsplittl-aux.simps*

**lemmas** *rsplit-aux-simps* = *bin-rsplit-aux-simps*

**lemmas** *th-if-simp1* = *split-if* [where  $P = op = l$ ,  
 $\text{THEN } \text{iffD1}, \text{ THEN } \text{conjunct1}, \text{ THEN } mp, \text{ standard}$ ]

**lemmas** *th-if-simp2* = *split-if* [where  $P = op = l$ ,  
 $\text{THEN } \text{iffD1}, \text{ THEN } \text{conjunct2}, \text{ THEN } mp, \text{ standard}$ ]

**lemmas** *rsplit-aux-simp1s* = *rsplit-aux-simps* [THEN *th-if-simp1*]

**lemmas** *rsplit-aux-simp2ls* = *rsplit-aux-simps* [THEN *th-if-simp2*]

**lemmas** *bin-rsplit-aux-simp2s* [simp] = *rsplit-aux-simp2ls* [unfolded *Let-def*]

**lemmas** *rsbscl* = *bin-rsplit-aux-simp2s* (2)

**lemmas** *rsplit-aux-0-simps* [simp] =

*rsplit-aux-simp1s* [OF *disjI1*] *rsplit-aux-simp1s* [OF *disjI2*]

**lemma** *bin-rsplit-aux-append*:

$\text{bin-rsplit-aux } n \ m \ c \ (\text{bs } @ \ \text{cs}) = \text{bin-rsplit-aux } n \ m \ c \ \text{bs } @ \ \text{cs}$   
 $\langle \text{proof} \rangle$

**lemma** *bin-rsplittl-aux-append*:

$\text{bin-rsplittl-aux } n \ m \ c \ (\text{bs } @ \ \text{cs}) = \text{bin-rsplittl-aux } n \ m \ c \ \text{bs } @ \ \text{cs}$   
 $\langle \text{proof} \rangle$

**lemmas** *rsplit-aux-apps* [where  $\text{bs} = []$ ] =

*bin-rsplit-aux-append* *bin-rsplittl-aux-append*

**lemmas** *rsplit-def-auxs* = *bin-rsplit-def* *bin-rsplittl-def*

**lemmas** *rsplit-aux-alts* = *rsplit-aux-apps*

[unfolded *append-Nil* *rsplit-def-auxs* [symmetric]]

**lemma** *bin-split-minus*:  $0 < n ==> \text{bin-split } (\text{Suc } (n - 1)) \ w = \text{bin-split } n \ w$

$\langle \text{proof} \rangle$

**lemmas** *bin-split-minus-simp* =  
*bin-split.Suc [THEN [2] bin-split-minus [symmetric, THEN trans], standard]*

**lemma** *bin-split-pred-simp [simp]*:  
 $(0::\text{nat}) < \text{number-of } \text{bin} \implies$   
 $\text{bin-split } (\text{number-of } \text{bin}) \ w =$   
 $(\text{let } (w1, w2) = \text{bin-split } (\text{number-of } (\text{Int.pred } \text{bin})) \ (\text{bin-rest } w)$   
 $\text{in } (w1, w2 \text{ BIT } \text{bin-last } w))$   
 $\langle \text{proof} \rangle$

**declare** *bin-split-pred-simp [simp]*

**lemma** *bin-rsplit-aux-simp-alt*:  
 $\text{bin-rsplit-aux } n \ m \ c \ bs =$   
 $(\text{if } m = 0 \ \vee \ n = 0$   
 $\text{then } bs$   
 $\text{else let } (a, b) = \text{bin-split } n \ c \text{ in } \text{bin-rsplit } n \ (m - n, a) \ @ \ b \ \# \ bs)$   
 $\langle \text{proof} \rangle$

**lemmas** *bin-rsplit-simp-alt* =  
 $\text{trans } [OF \ \text{bin-rsplit-def}$   
 $\text{bin-rsplit-aux-simp-alt, standard}]$

**lemmas** *bthrs* = *bin-rsplit-simp-alt [THEN [2] trans]*

**lemma** *bin-rsplit-size-sign' [rule-format]* :  
 $n > 0 \implies (\text{ALL } nw \ w. \text{rev } sw = \text{bin-rsplit } n \ (nw, w) \dashrightarrow$   
 $(\text{ALL } v: \text{set } sw. \text{bintrunc } n \ v = v))$   
 $\langle \text{proof} \rangle$

**lemmas** *bin-rsplit-size-sign* = *bin-rsplit-size-sign' [OF asm-rl*  
 $\text{rev-rev-ident [THEN trans] set-rev [THEN equalityD2 [THEN subsetD]],}$   
 $\text{standard}]$

**lemma** *bin-nth-rsplit [rule-format]* :  
 $n > 0 \implies m < n \implies (\text{ALL } w \ k \ nw. \text{rev } sw = \text{bin-rsplit } n \ (nw, w) \dashrightarrow$   
 $k < \text{size } sw \dashrightarrow \text{bin-nth } (sw ! k) \ m = \text{bin-nth } w \ (k * n + m))$   
 $\langle \text{proof} \rangle$

**lemma** *bin-rsplit-all*:  
 $0 < nw \implies nw \leq n \implies \text{bin-rsplit } n \ (nw, w) = [\text{bintrunc } n \ w]$   
 $\langle \text{proof} \rangle$

**lemma** *bin-rsplit-l [rule-format]* :  
 $\text{ALL } \text{bin}. \text{bin-rsplitl } n \ (m, \text{bin}) = \text{bin-rsplit } n \ (m, \text{bintrunc } m \ \text{bin})$   
 $\langle \text{proof} \rangle$

**lemma** *bin-rsplit-rcat* [rule-format] :

$n > 0 \dashrightarrow \text{bin-rsplit } n \ (n * \text{size } ws, \text{bin-rcat } n \ ws) = \text{map } (\text{bintrunc } n) \ ws$   
 ⟨proof⟩

**lemma** *bin-rsplit-aux-len-le* [rule-format] :

$\forall ws \ m. \ n \neq 0 \longrightarrow ws = \text{bin-rsplit-aux } n \ nw \ w \ bs \longrightarrow$   
 $\text{length } ws \leq m \iff nw + \text{length } bs * n \leq m * n$   
 ⟨proof⟩

**lemma** *bin-rsplit-len-le*:

$n \neq 0 \dashrightarrow ws = \text{bin-rsplit } n \ (nw, w) \dashrightarrow (\text{length } ws \leq m) = (nw \leq m * n)$   
 ⟨proof⟩

**lemma** *bin-rsplit-aux-len* [rule-format] :

$n \neq 0 \dashrightarrow \text{length } (\text{bin-rsplit-aux } n \ nw \ w \ cs) =$   
 $(nw + n - 1) \text{ div } n + \text{length } cs$   
 ⟨proof⟩

**lemma** *bin-rsplit-len*:

$n \neq 0 \implies \text{length } (\text{bin-rsplit } n \ (nw, w)) = (nw + n - 1) \text{ div } n$   
 ⟨proof⟩

**lemma** *bin-rsplit-aux-len-indep*:

$n \neq 0 \implies \text{length } bs = \text{length } cs \implies$   
 $\text{length } (\text{bin-rsplit-aux } n \ nw \ v \ bs) =$   
 $\text{length } (\text{bin-rsplit-aux } n \ nw \ w \ cs)$   
 ⟨proof⟩

**lemma** *bin-rsplit-len-indep*:

$n \neq 0 \implies \text{length } (\text{bin-rsplit } n \ (nw, v)) = \text{length } (\text{bin-rsplit } n \ (nw, w))$   
 ⟨proof⟩

end

## 8 TdThs: Type Definition Theorems

**theory** *TdThs*

**imports** *Main*

**begin**

## 9 More lemmas about normal type definitions

**lemma**

*tdD1*: *type-definition* *Rep Abs A*  $\implies \forall x. \text{Rep } x \in A$  **and**

*tdD2*: *type-definition* *Rep Abs A*  $\implies \forall x. \text{Abs } (\text{Rep } x) = x$  **and**

*tdD3*: *type-definition* *Rep Abs A*  $\implies \forall y. y \in A \longrightarrow \text{Rep } (\text{Abs } y) = y$   
 ⟨*proof*⟩

**lemma** *td-nat-int*:  
*type-definition* *int nat* (*Collect* (*op*  $\leq 0$ ))  
 ⟨*proof*⟩

**context** *type-definition*  
**begin**

**lemmas** *Rep'* [*iff*] = *Rep* [*simplified*]

**declare** *Rep-inverse* [*simp*] *Rep-inject* [*simp*]

**lemma** *Abs-eqD*: *Abs* *x* = *Abs* *y*  $\implies x \in A \implies y \in A \implies x = y$   
 ⟨*proof*⟩

**lemma** *Abs-inverse'*:  
*r* : *A*  $\implies \text{Abs } r = a \implies \text{Rep } a = r$   
 ⟨*proof*⟩

**lemma** *Rep-comp-inverse*:  
*Rep* *o* *f* = *g*  $\implies \text{Abs } o$  *g* = *f*  
 ⟨*proof*⟩

**lemma** *Rep-eqD* [*elim!*]: *Rep* *x* = *Rep* *y*  $\implies x = y$   
 ⟨*proof*⟩

**lemma** *Rep-inverse'*: *Rep* *a* = *r*  $\implies \text{Abs } r = a$   
 ⟨*proof*⟩

**lemma** *comp-Abs-inverse*:  
*f* *o* *Abs* = *g*  $\implies g$  *o* *Rep* = *f*  
 ⟨*proof*⟩

**lemma** *set-Rep*:  
*A* = *range* *Rep*  
 ⟨*proof*⟩

**lemma** *set-Rep-Abs*: *A* = *range* (*Rep* *o* *Abs*)  
 ⟨*proof*⟩

**lemma** *Abs-inj-on*: *inj-on* *Abs* *A*  
 ⟨*proof*⟩

**lemma** *image*: *Abs* ‘ *A* = *UNIV*  
 ⟨*proof*⟩

**lemmas** *td-thm* = *type-definition-axioms*

**lemma** *fns1*:  
 $Rep \circ fa = fr \circ Rep \mid fa \circ Abs = Abs \circ fr ==> Abs \circ fr \circ Rep = fa$   
 ⟨proof⟩

**lemmas** *fns1a* = *disjI1* [THEN *fns1*]  
**lemmas** *fns1b* = *disjI2* [THEN *fns1*]

**lemma** *fns4*:  
 $Rep \circ fa \circ Abs = fr ==>$   
 $Rep \circ fa = fr \circ Rep \ \& \ fa \circ Abs = Abs \circ fr$   
 ⟨proof⟩

**end**

**interpretation** *nat-int*: *type-definition* [*int nat Collect (op <= 0)*]  
 ⟨proof⟩  
**declare** *Nat.induct* [*case-names 0 Suc, induct type*]  
**declare** *Nat.exhaust* [*case-names 0 Suc, cases type*]

## 9.1 Extended form of type definition predicate

**lemma** *td-conds*:  
 $norm \circ norm = norm ==> (fr \circ norm = norm \circ fr) =$   
 $(norm \circ fr \circ norm = fr \circ norm \ \& \ norm \circ fr \circ norm = norm \circ fr)$   
 ⟨proof⟩

**lemma** *fn-comm-power*:  
 $fa \circ tr = tr \circ fr ==> fa \wedge n \circ tr = tr \circ fr \wedge n$   
 ⟨proof⟩

**lemmas** *fn-comm-power'* =  
*ext* [THEN *fn-comm-power*, THEN *fun-cong*, *unfolded o-def*, *standard*]

**locale** *td-ext* = *type-definition* +  
**fixes** *norm*  
**assumes** *eq-norm*:  $\bigwedge x. Rep (Abs x) = norm x$   
**begin**

**lemma** *Abs-norm* [*simp*]:  
 $Abs (norm x) = Abs x$   
 ⟨proof⟩

**lemma** *td-th*:  
 $g \circ Abs = f ==> f (Rep x) = g x$   
 ⟨proof⟩

**lemma** *eq-norm'*:  $Rep \circ Abs = norm$

$\langle proof \rangle$

**lemma** *norm-Rep [simp]*:  $norm (Rep x) = Rep x$   
 $\langle proof \rangle$

**lemmas** *td = td-thm*

**lemma** *set-iff-norm*:  $w : A \leftrightarrow w = norm w$   
 $\langle proof \rangle$

**lemma** *inverse-norm*:  
 $(Abs n = w) = (Rep w = norm n)$   
 $\langle proof \rangle$

**lemma** *norm-eq-iff*:  
 $(norm x = norm y) = (Abs x = Abs y)$   
 $\langle proof \rangle$

**lemma** *norm-comps*:  
 $Abs o norm = Abs$   
 $norm o Rep = Rep$   
 $norm o norm = norm$   
 $\langle proof \rangle$

**lemmas** *norm-norm [simp] = norm-comps*

**lemma** *fns5*:  
 $Rep o fa o Abs = fr ==>$   
 $fr o norm = fr \ \& \ norm o fr = fr$   
 $\langle proof \rangle$

**lemma** *fns2*:  
 $Abs o fr o Rep = fa ==>$   
 $(norm o fr o norm = fr o norm) = (Rep o fa = fr o Rep)$   
 $\langle proof \rangle$

**lemma** *fns3*:  
 $Abs o fr o Rep = fa ==>$   
 $(norm o fr o norm = norm o fr) = (fa o Abs = Abs o fr)$   
 $\langle proof \rangle$

**lemma** *fns*:  
 $fr o norm = norm o fr ==>$   
 $(fa o Abs = Abs o fr) = (Rep o fa = fr o Rep)$   
 $\langle proof \rangle$

**lemma** *range-norm*:  
 $range (Rep o Abs) = A$

```

    <proof>

end

lemmas td-ext-def' =
  td-ext-def [unfolded type-definition-def td-ext-axioms-def]

end

```

## 10 WordDefinition: Definition of Word Type

```

theory WordDefinition
imports Size BinBoolList TdThs
begin

typedef (open word) 'a word
  = {(0::int) ..< 2^len-of TYPE('a::len0)} <proof>

instantiation word :: (len0) size
begin

definition
  word-size: size (w :: 'a word) = len-of TYPE('a)

instance <proof>

end

definition
  — representation of words using unsigned or signed bins, only difference in these
  is the type class
  word-of-int :: int ⇒ 'a::len0 word
where
  [code func del]: word-of-int w = Abs-word (bintrunc (len-of TYPE ('a)) w)

code-datatype word-of-int

```

### 10.1 Type conversions and casting

```

constdefs
  — uint and sint cast a word to an integer, uint treats the word as unsigned, sint
  treats the most-significant-bit as a sign bit
  uint :: 'a :: len0 word => int
  uint w == Rep-word w
  sint :: 'a :: len word => int
  sint-uint: sint w == sbintrunc (len-of TYPE ('a) - 1) (uint w)
  unat :: 'a :: len0 word => nat

```

*unat w == nat (uint w)*

— the sets of integers representing the words

*uints :: nat => int set*

*uints n == range (bintrunc n)*

*sints :: nat => int set*

*sints n == range (sbintrunc (n - 1))*

*unats :: nat => nat set*

*unats n == {i. i < 2 ^ n}*

*norm-sint :: nat => int => int*

*norm-sint n w == (w + 2 ^ (n - 1)) mod 2 ^ n - 2 ^ (n - 1)*

— cast a word to a different length

*scast :: 'a :: len word => 'b :: len word*

*scast w == word-of-int (sint w)*

*ucast :: 'a :: len0 word => 'b :: len0 word*

*ucast w == word-of-int (uint w)*

— whether a cast (or other) function is to a longer or shorter length

*source-size :: ('a :: len0 word => 'b) => nat*

*source-size c == let arb = arbitrary ; x = c arb in size arb*

*target-size :: ('a => 'b :: len0 word) => nat*

*target-size c == size (c arbitrary)*

*is-up :: ('a :: len0 word => 'b :: len0 word) => bool*

*is-up c == source-size c <= target-size c*

*is-down :: ('a :: len0 word => 'b :: len0 word) => bool*

*is-down c == target-size c <= source-size c*

### constdefs

*of-bl :: bool list => 'a :: len0 word*

*of-bl bl == word-of-int (bl-to-bin bl)*

*to-bl :: 'a :: len0 word => bool list*

*to-bl w ==*

*bin-to-bl (len-of TYPE ('a)) (uint w)*

*word-reverse :: 'a :: len0 word => 'a word*

*word-reverse w == of-bl (rev (to-bl w))*

### constdefs

*word-int-case :: (int => 'b) => ('a :: len0 word) => 'b*

*word-int-case f w == f (uint w)*

### syntax

*of-int :: int => 'a*

### translations

*case x of of-int y => b == word-int-case (%y. b) x*



## 10.2 Arithmetic operations

**declare** *uint-def* [*code func del*]

**lemma** [*code func*]: *uint* (*word-of-int* *w* :: '*a*::*len0* *word*) = *bintrunc* (*len-of TYPE*('a'))  
*w*  
 ⟨*proof*⟩

**instantiation** *word* :: (*len0*) {*number*, *uminus*, *minus*, *plus*, *one*, *zero*, *times*,  
*Divides.div*, *power*, *ord*, *bit*}

**begin**

**definition**

*word-0-wi*:  $0 = \text{word-of-int } 0$

**definition**

*word-1-wi*:  $1 = \text{word-of-int } 1$

**definition**

*word-add-def*:  $a + b = \text{word-of-int } (\text{uint } a + \text{uint } b)$

**definition**

*word-sub-wi*:  $a - b = \text{word-of-int } (\text{uint } a - \text{uint } b)$

**definition**

*word-minus-def*:  $- a = \text{word-of-int } (- \text{uint } a)$

**definition**

*word-mult-def*:  $a * b = \text{word-of-int } (\text{uint } a * \text{uint } b)$

**definition**

*word-div-def*:  $a \text{ div } b = \text{word-of-int } (\text{uint } a \text{ div } \text{uint } b)$

**definition**

*word-mod-def*:  $a \text{ mod } b = \text{word-of-int } (\text{uint } a \text{ mod } \text{uint } b)$

**primrec** *power-word* **where**

$(a::'a \text{ word}) \wedge 0 = 1$   
 $| (a::'a \text{ word}) \wedge \text{Suc } n = a * a \wedge n$

**definition**

*word-number-of-def*:  $\text{number-of } w = \text{word-of-int } w$

**definition**

*word-le-def*:  $a \leq b \longleftrightarrow \text{uint } a \leq \text{uint } b$

**definition**

*word-less-def*:  $x < y \longleftrightarrow x \leq y \wedge x \neq (y :: 'a \text{ word})$

**definition**

*word-and-def:*  
 $(a :: 'a \text{ word}) \text{ AND } b = \text{word-of-int } (\text{uint } a \text{ AND } \text{uint } b)$

**definition**

*word-or-def:*  
 $(a :: 'a \text{ word}) \text{ OR } b = \text{word-of-int } (\text{uint } a \text{ OR } \text{uint } b)$

**definition**

*word-xor-def:*  
 $(a :: 'a \text{ word}) \text{ XOR } b = \text{word-of-int } (\text{uint } a \text{ XOR } \text{uint } b)$

**definition**

*word-not-def:*  
 $\text{NOT } (a :: 'a \text{ word}) = \text{word-of-int } (\text{NOT } (\text{uint } a))$

**instance**  $\langle \text{proof} \rangle$

**end**

**definition**

$\text{word-succ} :: 'a :: \text{len0 word} \Rightarrow 'a \text{ word}$

**where**

$\text{word-succ } a = \text{word-of-int } (\text{Int.succ } (\text{uint } a))$

**definition**

$\text{word-pred} :: 'a :: \text{len0 word} \Rightarrow 'a \text{ word}$

**where**

$\text{word-pred } a = \text{word-of-int } (\text{Int.pred } (\text{uint } a))$

**constdefs**

$\text{udvd} :: 'a :: \text{len word} \Rightarrow 'a :: \text{len word} \Rightarrow \text{bool}$  (**infixl**  $\text{udvd}$  50)  
 $a \text{ udvd } b == \text{EX } n \geq 0. \text{uint } b = n * \text{uint } a$

$\text{word-sle} :: 'a :: \text{len word} \Rightarrow 'a \text{ word} \Rightarrow \text{bool}$   $((-/ \leq s -) [50, 51] 50)$   
 $a \leq s b == \text{sint } a \leq \text{sint } b$

$\text{word-sless} :: 'a :: \text{len word} \Rightarrow 'a \text{ word} \Rightarrow \text{bool}$   $((-/ < s -) [50, 51] 50)$   
 $(x < s y) == (x \leq s y \ \& \ x \sim = y)$

**10.3 Bit-wise operations**

**instantiation**  $\text{word} :: (\text{len0}) \text{ bits}$

**begin**

**definition**

*word-test-bit-def:*  $\text{test-bit } a = \text{bin-nth } (\text{uint } a)$

**definition**

*word-set-bit-def:*  $\text{set-bit } a \ n \ x =$

*word-of-int* (*bin-sc* *n* (*If* *x* *bit.B1* *bit.B0*) (*uint* *a*))

**definition**

*word-set-bits-def*: (*BITS* *n*. *f* *n*) = *of-bl* (*bl-of-nth* (*len-of TYPE* (*'a*)) *f*)

**definition**

*word-lsb-def*: *lsb* *a*  $\longleftrightarrow$  *bin-last* (*uint* *a*) = *bit.B1*

**definition** *shiftrl1* :: *'a* *word*  $\Rightarrow$  *'a* *word* **where**

*shiftrl1* *w* = *word-of-int* (*uint* *w* *BIT* *bit.B0*)

**definition** *shiftr1* :: *'a* *word*  $\Rightarrow$  *'a* *word* **where**

— shift right as unsigned or as signed, ie logical or arithmetic

*shiftr1* *w* = *word-of-int* (*bin-rest* (*uint* *w*))

**definition**

*shiftrl-def*: *w* << *n* = (*shiftrl1* ^ *n*) *w*

**definition**

*shiftr-def*: *w* >> *n* = (*shiftr1* ^ *n*) *w*

**instance**  $\langle proof \rangle$

**end**

**instantiation** *word* :: (*len*) *bitss*

**begin**

**definition**

*word-msb-def*:

*msb* *a*  $\longleftrightarrow$  *bin-sign* (*sint* *a*) = *Int.Min*

**instance**  $\langle proof \rangle$

**end**

**constdefs**

*setBit* :: *'a* :: *len0* *word*  $\Rightarrow$  *nat*  $\Rightarrow$  *'a* *word*

*setBit* *w* *n* == *set-bit* *w* *n* *True*

*clearBit* :: *'a* :: *len0* *word*  $\Rightarrow$  *nat*  $\Rightarrow$  *'a* *word*

*clearBit* *w* *n* == *set-bit* *w* *n* *False*

## 10.4 Shift operations

**constdefs**

*sshiftr1* :: *'a* :: *len* *word*  $\Rightarrow$  *'a* *word*

*sshiftr1* *w* == *word-of-int* (*bin-rest* (*sint* *w*))

```

bshiftr1 :: bool => 'a :: len word => 'a word
bshiftr1 b w == of-bl (b # butlast (to-bl w))

sshiftr :: 'a :: len word => nat => 'a word (infixl >>> 55)
w >>> n == (sshiftr1 ^ n) w

mask :: nat => 'a :: len word
mask n == (1 << n) - 1

revcast :: 'a :: len0 word => 'b :: len0 word
revcast w == of-bl (takefill False (len-of TYPE('b)) (to-bl w))

slice1 :: nat => 'a :: len0 word => 'b :: len0 word
slice1 n w == of-bl (takefill False n (to-bl w))

slice :: nat => 'a :: len0 word => 'b :: len0 word
slice n w == slice1 (size w - n) w

```

## 10.5 Rotation

### constdefs

```

rotater1 :: 'a list => 'a list
rotater1 ys ==
  case ys of [] => [] | x # xs => last ys # butlast ys

rotater :: nat => 'a list => 'a list
rotater n == rotater1 ^ n

word-rotr :: nat => 'a :: len0 word => 'a :: len0 word
word-rotr n w == of-bl (rotater n (to-bl w))

word-rotl :: nat => 'a :: len0 word => 'a :: len0 word
word-rotl n w == of-bl (rotate n (to-bl w))

word-roti :: int => 'a :: len0 word => 'a :: len0 word
word-roti i w == if i >= 0 then word-rotr (nat i) w
                  else word-rotl (nat (- i)) w

```

## 10.6 Split and cat operations

### constdefs

```

word-cat :: 'a :: len0 word => 'b :: len0 word => 'c :: len0 word
word-cat a b == word-of-int (bin-cat (uint a) (len-of TYPE('b)) (uint b))

word-split :: 'a :: len0 word => ('b :: len0 word) * ('c :: len0 word)
word-split a ==
  case bin-split (len-of TYPE('c)) (uint a) of
    (u, v) => (word-of-int u, word-of-int v)

word-rcat :: 'a :: len0 word list => 'b :: len0 word

```

```

word-rcat ws ==
word-of-int (bin-rcat (len-of TYPE ('a)) (map uint ws))

word-rsplit :: 'a :: len0 word => 'b :: len word list
word-rsplit w ==
map word-of-int (bin-rsplit (len-of TYPE ('b)) (len-of TYPE ('a), uint w))

```

**constdefs**

— Largest representable machine integer.

```

max-word :: 'a::len word
max-word ≡ word-of-int (2len-of TYPE('a) - 1)

```

**consts**

```

of-bool :: bool ⇒ 'a::len word

```

**primrec**

```

of-bool False = 0
of-bool True = 1

```

**lemmas** *of-nth-def* = *word-set-bits-def*

**lemmas** *word-size-gt-0 [iff]* =  
*xtr1 [OF word-size len-gt-0, standard]*  
**lemmas** *lens-gt-0* = *word-size-gt-0 len-gt-0*  
**lemmas** *lens-not-0 [iff]* = *lens-gt-0 [THEN gr-implies-not0, standard]*

**lemma** *uints-num*: *uints n* = {*i*. 0 ≤ *i* ∧ *i* < 2<sup>*n*</sup>}

⟨*proof*⟩

**lemma** *sints-num*: *sints n* = {*i*. −(2<sup>*n* − 1</sup>) ≤ *i* ∧ *i* < 2<sup>*n*</sup>}

⟨*proof*⟩

**lemmas** *atLeastLessThan-alt* = *atLeastLessThan-def [unfolded*  
*atLeast-def lessThan-def Collect-conj-eq [symmetric]]*

**lemma** *mod-in-reps*: *m* > 0 ==> *y mod m* : {0::int ..< *m*}

⟨*proof*⟩

**lemma**

*Rep-word-0:0* ≤ *Rep-word x* **and**  
*Rep-word-lt*: *Rep-word (x::'a::len0 word)* < 2<sup>len-of TYPE('a)</sup>

⟨*proof*⟩

**lemma** *Rep-word-mod-same*:

*Rep-word x mod 2<sup>len-of TYPE('a)</sup>* = *Rep-word (x::'a::len0 word)*

⟨*proof*⟩

**lemma** *td-ext-uint*:

*td-ext (uint :: 'a word => int) word-of-int (uints (len-of TYPE('a::len0)))*

(%w::int. w mod 2 ^ len-of TYPE('a))  
 <proof>

**lemmas** int-word-uint = td-ext-uint [THEN td-ext.eq-norm, standard]

**interpretation** word-uint:  
 td-ext [uint::'a::len0 word  $\Rightarrow$  int  
   word-of-int  
   uints (len-of TYPE('a::len0))  
    $\lambda w. w \text{ mod } 2 \wedge \text{len-of TYPE('a::len0)}$ ]  
 <proof>

**lemmas** td-uint = word-uint.td-thm

**lemmas** td-ext-ubin = td-ext-uint  
 [simplified len-gt-0 no-bintr-alt1 [symmetric]]

**interpretation** word-ubin:  
 td-ext [uint::'a::len0 word  $\Rightarrow$  int  
   word-of-int  
   uints (len-of TYPE('a::len0))  
   bintrunc (len-of TYPE('a::len0))]  
 <proof>

**lemma** sint-sbintrunc':  
 sint (word-of-int bin :: 'a word) =  
 (sbintrunc (len-of TYPE ('a :: len) - 1) bin)  
 <proof>

**lemma** uint-sint:  
 uint w = bintrunc (len-of TYPE('a)) (sint (w :: 'a :: len word))  
 <proof>

**lemma** bintr-uint':  
 $n \geq \text{size } w \Rightarrow \text{bintrunc } n \text{ (uint } w) = \text{uint } w$   
 <proof>

**lemma** wi-bintr':  
 $wb = \text{word-of-int bin} \Rightarrow n \geq \text{size } wb \Rightarrow$   
 $\text{word-of-int (bintrunc } n \text{ bin)} = wb$   
 <proof>

**lemmas** bintr-uint = bintr-uint' [unfolded word-size]

**lemmas** wi-bintr = wi-bintr' [unfolded word-size]

**lemma** td-ext-sbin:  
 td-ext (sint :: 'a word  $\Rightarrow$  int) word-of-int (uints (len-of TYPE('a::len)))  
 (sbintrunc (len-of TYPE('a) - 1))  
 <proof>

**lemmas** *td-ext-sint = td-ext-sbin*  
*[simplified len-gt-0 no-sbintr-alt2 Suc-pred' [symmetric]]*

**interpretation** *word-sint:*  
*td-ext [sint :: 'a::len word => int*  
*word-of-int*  
*sints (len-of TYPE('a::len))*  
*%w. (w + 2<sup>(len-of TYPE('a::len) - 1)</sup>) mod 2<sup>len-of TYPE('a::len) -</sup>*  
*2<sup>(len-of TYPE('a::len) - 1)</sup>]*  
*<proof>*

**interpretation** *word-sbin:*  
*td-ext [sint :: 'a::len word => int*  
*word-of-int*  
*sints (len-of TYPE('a::len))*  
*sbintrunc (len-of TYPE('a::len) - 1)]*  
*<proof>*

**lemmas** *int-word-sint = td-ext-sint [THEN td-ext.eq-norm, standard]*

**lemmas** *td-sint = word-sint.td*

**lemma** *word-number-of-alt: number-of b == word-of-int (number-of b)*  
*<proof>*

**lemma** *word-no-wi: number-of = word-of-int*  
*<proof>*

**lemma** *to-bl-def':*  
*(to-bl :: 'a :: len0 word => bool list) =*  
*bin-to-bl (len-of TYPE('a)) o uint*  
*<proof>*

**lemmas** *word-reverse-no-def [simp] = word-reverse-def [of number-of w, standard]*

**lemmas** *uints-mod = uints-def [unfolded no-bintr-alt1]*

**lemma** *uint-bintrunc: uint (number-of bin :: 'a word) =*  
*number-of (bintrunc (len-of TYPE ('a :: len0)) bin)*  
*<proof>*

**lemma** *sint-sbintrunc: sint (number-of bin :: 'a word) =*  
*number-of (sbintrunc (len-of TYPE ('a :: len) - 1) bin)*  
*<proof>*

**lemma** *unat-bintrunc:*  
*unat (number-of bin :: 'a :: len0 word) =*

*number-of* (*bintrunc* (*len-of TYPE*('a)) *bin*)  
 ⟨*proof*⟩

**declare**

*uint-bintrunc* [*simp*]  
*sint-sbintrunc* [*simp*]  
*unat-bintrunc* [*simp*]

**lemma** *size-0-eq*: *size* (*w* :: 'a :: len0 word) = 0 ==> *v* = *w*  
 ⟨*proof*⟩

**lemmas** *uint-lem* = *word-uint.Rep* [*unfolded uints-num mem-Collect-eq*]  
**lemmas** *sint-lem* = *word-sint.Rep* [*unfolded sints-num mem-Collect-eq*]  
**lemmas** *uint-ge-0* [*iff*] = *uint-lem* [*THEN conjunct1, standard*]  
**lemmas** *uint-lt2p* [*iff*] = *uint-lem* [*THEN conjunct2, standard*]  
**lemmas** *sint-ge* = *sint-lem* [*THEN conjunct1, standard*]  
**lemmas** *sint-lt* = *sint-lem* [*THEN conjunct2, standard*]

**lemma** *sign-uint-Pls* [*simp*]:  
*bin-sign* (*uint x*) = *Int.Pl*s  
 ⟨*proof*⟩

**lemmas** *uint-m2p-neg* = *iffD2* [*OF diff-less-0-iff-less uint-lt2p, standard*]  
**lemmas** *uint-m2p-not-non-neg* =  
*iffD2* [*OF linorder-not-le uint-m2p-neg, standard*]

**lemma** *lt2p-lem*:  
*len-of TYPE*('a) <= *n* ==> *uint* (*w* :: 'a :: len0 word) < 2 ^ *n*  
 ⟨*proof*⟩

**lemmas** *uint-le-0-iff* [*simp*] =  
*uint-ge-0* [*THEN leD, THEN linorder-antisym-conv1, standard*]

**lemma** *uint-nat*: *uint w* == *int* (*unat w*)  
 ⟨*proof*⟩

**lemma** *uint-number-of*:  
*uint* (*number-of b* :: 'a :: len0 word) = *number-of b mod 2 ^ len-of TYPE*('a)  
 ⟨*proof*⟩

**lemma** *unat-number-of*:  
*bin-sign b* = *Int.Pl*s ==>  
*unat* (*number-of b*::'a::len0 word) = *number-of b mod 2 ^ len-of TYPE* ('a)  
 ⟨*proof*⟩

**lemma** *sint-number-of*: *sint* (*number-of b* :: 'a :: len word) = (*number-of b* +  
 2 ^ (*len-of TYPE*('a) - 1)) mod 2 ^ *len-of TYPE*('a) -  
 2 ^ (*len-of TYPE*('a) - 1)



$\langle \text{proof} \rangle$

**lemma** *word-of-int-bin* [simp] :  
 $(\text{word-of-int } (\text{number-of bin}) :: 'a :: \text{len0 word}) = (\text{number-of bin})$   
 $\langle \text{proof} \rangle$

**lemma** *word-int-case-wi*:  
 $\text{word-int-case } f \ (\text{word-of-int } i :: 'b \text{ word}) =$   
 $f \ (i \bmod 2 \wedge \text{len-of TYPE('b::len0)})$   
 $\langle \text{proof} \rangle$

**lemma** *word-int-split*:  
 $P \ (\text{word-int-case } f \ x) =$   
 $(\text{ALL } i. \ x = (\text{word-of-int } i :: 'b :: \text{len0 word}) \ \&$   
 $0 \leq i \ \& \ i < 2 \wedge \text{len-of TYPE('b)} \longrightarrow P \ (f \ i))$   
 $\langle \text{proof} \rangle$

**lemma** *word-int-split-asm*:  
 $P \ (\text{word-int-case } f \ x) =$   
 $(\sim (\text{EX } n. \ x = (\text{word-of-int } n :: 'b::\text{len0 word}) \ \&$   
 $0 \leq n \ \& \ n < 2 \wedge \text{len-of TYPE('b::len0)} \ \& \sim P \ (f \ n)))$   
 $\langle \text{proof} \rangle$

**lemmas** *uint-range'* =  
 $\text{word-uint.Rep} \ [\text{unfolded uints-num mem-Collect-eq, standard}]$   
**lemmas** *sint-range'* =  $\text{word-sint.Rep} \ [\text{unfolded One-nat-def}$   
 $\text{sints-num mem-Collect-eq, standard}]$

**lemma** *uint-range-size*:  $0 \leq \text{uint } w \ \& \ \text{uint } w < 2 \wedge \text{size } w$   
 $\langle \text{proof} \rangle$

**lemma** *sint-range-size*:  
 $-(2 \wedge (\text{size } w - \text{Suc } 0)) \leq \text{sint } w \ \& \ \text{sint } w < 2 \wedge (\text{size } w - \text{Suc } 0)$   
 $\langle \text{proof} \rangle$

**lemmas** *sint-above-size* = *sint-range-size*  
 $[\text{THEN conjunct2, THEN } [2] \text{ xtr8, folded One-nat-def, standard}]$

**lemmas** *sint-below-size* = *sint-range-size*  
 $[\text{THEN conjunct1, THEN } [2] \text{ order-trans, folded One-nat-def, standard}]$

**lemma** *test-bit-eq-iff*:  $(\text{test-bit } (u::'a::\text{len0 word}) = \text{test-bit } v) = (u = v)$   
 $\langle \text{proof} \rangle$

**lemma** *test-bit-size* [rule-format] :  $(w::'a::\text{len0 word}) !! n \longrightarrow n < \text{size } w$   
 $\langle \text{proof} \rangle$

**lemma** *word-eqI* [rule-format] :  
**fixes**  $u :: 'a::\text{len0 word}$

**shows** (*ALL*  $n$ .  $n < \text{size } u \dashrightarrow u !! n = v !! n$ )  $\implies u = v$   
 ⟨*proof*⟩

**lemmas**  $\text{word-eqD} = \text{test-bit-eq-iff}$  [*THEN iffD2, THEN fun-cong, standard*]

**lemma**  $\text{test-bit-bin}'$ :  $w !! n = (n < \text{size } w \ \& \ \text{bin-nth} (\text{uint } w) \ n)$   
 ⟨*proof*⟩

**lemmas**  $\text{test-bit-bin} = \text{test-bit-bin}'$  [*unfolded word-size*]

**lemma**  $\text{bin-nth-uint-imp}'$ :  $\text{bin-nth} (\text{uint } w) \ n \dashrightarrow n < \text{size } w$   
 ⟨*proof*⟩

**lemma**  $\text{bin-nth-sint}'$ :  
 $n \geq \text{size } w \dashrightarrow \text{bin-nth} (\text{sint } w) \ n = \text{bin-nth} (\text{sint } w) (\text{size } w - 1)$   
 ⟨*proof*⟩

**lemmas**  $\text{bin-nth-uint-imp} = \text{bin-nth-uint-imp}'$  [*rule-format, unfolded word-size*]

**lemmas**  $\text{bin-nth-sint} = \text{bin-nth-sint}'$  [*rule-format, unfolded word-size*]

**lemma**  $\text{td-bl}$ :  
*type-definition* ( $\text{to-bl} :: 'a::\text{len0} \ \text{word} \Rightarrow \text{bool list}$ )  
*of-bl*  
 $\{\text{bl}. \text{length } \text{bl} = \text{len-of } \text{TYPE}('a)\}$   
 ⟨*proof*⟩

**interpretation**  $\text{word-bl}$ :  
*type-definition* [ $\text{to-bl} :: 'a::\text{len0} \ \text{word} \Rightarrow \text{bool list}$   
*of-bl*  
 $\{\text{bl}. \text{length } \text{bl} = \text{len-of } \text{TYPE}('a::\text{len0})\}$ ]  
 ⟨*proof*⟩

**lemma**  $\text{word-size-bl}$ :  $\text{size } w == \text{size} (\text{to-bl } w)$   
 ⟨*proof*⟩

**lemma**  $\text{to-bl-use-of-bl}$ :  
 $(\text{to-bl } w = \text{bl}) = (w = \text{of-bl } \text{bl} \ \wedge \ \text{length } \text{bl} = \text{length} (\text{to-bl } w))$   
 ⟨*proof*⟩

**lemma**  $\text{to-bl-word-rev}$ :  $\text{to-bl} (\text{word-reverse } w) = \text{rev} (\text{to-bl } w)$   
 ⟨*proof*⟩

**lemma**  $\text{word-rev-rev}$  [*simp*] :  $\text{word-reverse} (\text{word-reverse } w) = w$   
 ⟨*proof*⟩

**lemma**  $\text{word-rev-gal}$ :  $\text{word-reverse } w = u \implies \text{word-reverse } u = w$   
 ⟨*proof*⟩

**lemmas** *word-rev-gal'* = *sym* [THEN *word-rev-gal*, *symmetric*, *standard*]

**lemmas** *length-bl-gt-0* [iff] = *xtr1* [OF *word-bl.Rep'* *len-gt-0*, *standard*]

**lemmas** *bl-not-Nil* [iff] =

*length-bl-gt-0* [THEN *length-greater-0-conv* [THEN *iffD1*], *standard*]

**lemmas** *length-bl-neq-0* [iff] = *length-bl-gt-0* [THEN *gr-implies-not0*]

**lemma** *hd-bl-sign-sint*: *hd* (*to-bl w*) = (*bin-sign* (*sint w*) = *Int.Min*)  
 ⟨*proof*⟩

**lemma** *of-bl-drop'*:

*lend* = *length bl* − *len-of TYPE* ('*a* :: *len0*) ==>

*of-bl* (*drop lend bl*) = (*of-bl bl* :: '*a word*)

⟨*proof*⟩

**lemmas** *of-bl-no* = *of-bl-def* [*folded word-number-of-def*]

**lemma** *test-bit-of-bl*:

(*of-bl bl* :: '*a* :: *len0 word*) !! *n* = (*rev bl* ! *n* ∧ *n* < *len-of TYPE*('*a*) ∧ *n* < *length bl*)

⟨*proof*⟩

**lemma** *no-of-bl*:

(*number-of bin* :: '*a* :: *len0 word*) = *of-bl* (*bin-to-bl* (*len-of TYPE*('*a*)) *bin*)

⟨*proof*⟩

**lemma** *uint-bl*: *to-bl w* == *bin-to-bl* (*size w*) (*uint w*)

⟨*proof*⟩

**lemma** *to-bl-bin*: *bl-to-bin* (*to-bl w*) = *uint w*

⟨*proof*⟩

**lemma** *to-bl-of-bin*:

*to-bl* (*word-of-int bin* :: '*a* :: *len0 word*) = *bin-to-bl* (*len-of TYPE*('*a*)) *bin*

⟨*proof*⟩

**lemmas** *to-bl-no-bin* [*simp*] = *to-bl-of-bin* [*folded word-number-of-def*]

**lemma** *to-bl-to-bin* [*simp*] : *bl-to-bin* (*to-bl w*) = *uint w*

⟨*proof*⟩

**lemmas** *uint-bl-bin* [*simp*] = *trans* [OF *bin-bl-bin word-ubin.norm-Rep*, *standard*]

**lemmas** *num-AB-u* [*simp*] = *word-uint.Rep-inverse*

[*unfolded o-def word-number-of-def* [*symmetric*], *standard*]

**lemmas** *num-AB-s* [*simp*] = *word-sint.Rep-inverse*

[*unfolded o-def word-number-of-def* [*symmetric*], *standard*]

**lemma** *uints-unats*: *uints n = int ‘ unats n*  
 ⟨*proof*⟩

**lemma** *unats-uints*: *unats n = nat ‘ uints n*  
 ⟨*proof*⟩

**lemmas** *bintr-num = word-ubin.norm-eq-iff*  
 [*symmetric, folded word-number-of-def, standard*]

**lemmas** *sbintr-num = word-sbin.norm-eq-iff*  
 [*symmetric, folded word-number-of-def, standard*]

**lemmas** *num-of-bintr = word-ubin.Abs-norm* [*folded word-number-of-def, standard*]

**lemmas** *num-of-sbintr = word-sbin.Abs-norm* [*folded word-number-of-def, standard*]

**lemma** *num-of-bintr'*:  
*bintrunc (len-of TYPE('a :: len0)) a = b ==>*  
*number-of a = (number-of b :: 'a word)*  
 ⟨*proof*⟩

**lemma** *num-of-sbintr'*:  
*sbintrunc (len-of TYPE('a :: len) - 1) a = b ==>*  
*number-of a = (number-of b :: 'a word)*  
 ⟨*proof*⟩

**lemmas** *num-abs-bintr = sym [THEN trans,*  
*OF num-of-bintr word-number-of-def, standard]*

**lemmas** *num-abs-sbintr = sym [THEN trans,*  
*OF num-of-sbintr word-number-of-def, standard]*

**lemma** *ucast-id*: *ucast w = w*  
 ⟨*proof*⟩

**lemma** *scast-id*: *scast w = w*  
 ⟨*proof*⟩

**lemma** *ucast-bl*: *ucast w == of-bl (to-bl w)*  
 ⟨*proof*⟩

**lemma** *nth-ucast*:  
*(ucast w :: 'a :: len0 word) !! n = (w !! n & n < len-of TYPE('a))*  
 ⟨*proof*⟩

**lemma** *ucast-bintr* [*simp*]:  
 $ucast\ (number-of\ w :: 'a::len0\ word) =$   
 $number-of\ (bintrunc\ (len-of\ TYPE('a))\ w)$   
 ⟨*proof*⟩

**lemma** *scast-sbintr* [*simp*]:  
 $scast\ (number-of\ w :: 'a::len\ word) =$   
 $number-of\ (sbintrunc\ (len-of\ TYPE('a) - Suc\ 0)\ w)$   
 ⟨*proof*⟩

**lemmas** *source-size = source-size-def* [*unfolded Let-def word-size*]

**lemmas** *target-size = target-size-def* [*unfolded Let-def word-size*]

**lemmas** *is-down = is-down-def* [*unfolded source-size target-size*]

**lemmas** *is-up = is-up-def* [*unfolded source-size target-size*]

**lemmas** *is-up-down =*  
 $trans\ [OF\ is-up\ [THEN\ meta-eq-to-obj-eq]$   
 $is-down\ [THEN\ meta-eq-to-obj-eq,\ symmetric],$   
 $standard]$

**lemma** *down-cast-same'*:  $uc = ucast ==> is-down\ uc ==> uc = scast$   
 ⟨*proof*⟩

**lemma** *word-rev-tf'*:  
 $r = to-bl\ (of-bl\ bl) ==> r = rev\ (takefill\ False\ (length\ r)\ (rev\ bl))$   
 ⟨*proof*⟩

**lemmas** *word-rev-tf = refl* [*THEN word-rev-tf', unfolded word-bl.Rep', standard*]

**lemmas** *word-rep-drop = word-rev-tf* [*simplified takefill-alt,*  
*simplified, simplified rev-take, simplified*]

**lemma** *to-bl-ucast*:  
 $to-bl\ (ucast\ (w :: 'b::len0\ word) :: 'a::len0\ word) =$   
 $replicate\ (len-of\ TYPE('a) - len-of\ TYPE('b))\ False\ @$   
 $drop\ (len-of\ TYPE('b) - len-of\ TYPE('a))\ (to-bl\ w)$   
 ⟨*proof*⟩

**lemma** *ucast-up-app'*:  
 $uc = ucast ==> source-size\ uc + n = target-size\ uc ==>$   
 $to-bl\ (uc\ w) = replicate\ n\ False\ @ (to-bl\ w)$   
 ⟨*proof*⟩

**lemma** *ucast-down-drop'*:  
 $uc = ucast ==> source-size\ uc = target-size\ uc + n ==>$   
 $to-bl\ (uc\ w) = drop\ n\ (to-bl\ w)$   
 ⟨*proof*⟩

**lemma** *scast-down-drop'*:

$sc = scast ==> source\text{-}size\ sc = target\text{-}size\ sc + n ==>$   
 $to\text{-}bl\ (sc\ w) = drop\ n\ (to\text{-}bl\ w)$   
 ⟨proof⟩

**lemma** *sint-up-scast'*:

$sc = scast ==> is\text{-}up\ sc ==> sint\ (sc\ w) = sint\ w$   
 ⟨proof⟩

**lemma** *uint-up-ucast'*:

$uc = ucast ==> is\text{-}up\ uc ==> uint\ (uc\ w) = uint\ w$   
 ⟨proof⟩

**lemmas** *down-cast-same* = refl [THEN down-cast-same']

**lemmas** *ucast-up-app* = refl [THEN ucast-up-app']

**lemmas** *ucast-down-drop* = refl [THEN ucast-down-drop']

**lemmas** *scast-down-drop* = refl [THEN scast-down-drop']

**lemmas** *uint-up-ucast* = refl [THEN uint-up-ucast']

**lemmas** *sint-up-scast* = refl [THEN sint-up-scast']

**lemma** *ucast-up-ucast'*:  $uc = ucast ==> is\text{-}up\ uc ==> ucast\ (uc\ w) = ucast\ w$   
 ⟨proof⟩

**lemma** *scast-up-scast'*:  $sc = scast ==> is\text{-}up\ sc ==> scast\ (sc\ w) = scast\ w$   
 ⟨proof⟩

**lemma** *ucast-of-bl-up'*:

$w = of\text{-}bl\ bl ==> size\ bl \leq size\ w ==> ucast\ w = of\text{-}bl\ bl$   
 ⟨proof⟩

**lemmas** *ucast-up-ucast* = refl [THEN ucast-up-ucast']

**lemmas** *scast-up-scast* = refl [THEN scast-up-scast']

**lemmas** *ucast-of-bl-up* = refl [THEN ucast-of-bl-up']

**lemmas** *ucast-up-ucast-id* = trans [OF ucast-up-ucast ucast-id]

**lemmas** *scast-up-scast-id* = trans [OF scast-up-scast scast-id]

**lemmas** *isduu* = is-up-down [where  $c = ucast$ , THEN iffD2]

**lemmas** *isdus* = is-up-down [where  $c = scast$ , THEN iffD2]

**lemmas** *ucast-down-ucast-id* = isduu [THEN ucast-up-ucast-id]

**lemmas** *scast-down-scast-id* = isdus [THEN ucast-up-ucast-id]

**lemma** *up-ucast-surj*:

$is\text{-}up\ (ucast :: 'b::len0\ word ==> 'a::len0\ word) ==>$   
 $surj\ (ucast :: 'a\ word ==> 'b\ word)$   
 ⟨proof⟩

**lemma** *up-scast-surj*:

$is\text{-}up\ (scast :: 'b::len\ word ==> 'a::len\ word) ==>$

*surj* (*scast* :: 'a word => 'b word)  
 ⟨*proof*⟩

**lemma** *down-scast-inj*:

*is-down* (*scast* :: 'b::len word => 'a::len word) ==>  
*inj-on* (*ucast* :: 'a word => 'b word) *A*  
 ⟨*proof*⟩

**lemma** *down-ucast-inj*:

*is-down* (*ucast* :: 'b::len0 word => 'a::len0 word) ==>  
*inj-on* (*ucast* :: 'a word => 'b word) *A*  
 ⟨*proof*⟩

**lemma** *of-bl-append-same*: *of-bl* (*X* @ *to-bl w*) = *w*  
 ⟨*proof*⟩

**lemma** *ucast-down-no'*:

*uc* = *ucast* ==> *is-down uc* ==> *uc* (*number-of bin*) = *number-of bin*  
 ⟨*proof*⟩

**lemmas** *ucast-down-no* = *ucast-down-no'* [*OF refl*]

**lemma** *ucast-down-bl'*: *uc* = *ucast* ==> *is-down uc* ==> *uc* (*of-bl bl*) = *of-bl bl*  
 ⟨*proof*⟩

**lemmas** *ucast-down-bl* = *ucast-down-bl'* [*OF refl*]

**lemmas** *slice-def'* = *slice-def* [*unfolded word-size*]

**lemmas** *test-bit-def'* = *word-test-bit-def* [*THEN fun-cong*]

**lemmas** *word-log-defs* = *word-and-def word-or-def word-xor-def word-not-def*

**lemmas** *word-log-bin-defs* = *word-log-defs*

**end**

## 11 WordArith: Word Arithmetic

**theory** *WordArith*

**imports** *WordDefinition*

**begin**

**lemma** *word-less-alt*: (*a* < *b*) = (*uint a* < *uint b*)  
 ⟨*proof*⟩

**lemma** *signed-linorder*: *linorder word-sle word-sless*  
 ⟨*proof*⟩

**interpretation** *signed*: *linorder* [*word-sle word-sless*]

$\langle proof \rangle$

**lemmas** *word-arith-wis* =  
*word-add-def word-mult-def word-minus-def*  
*word-succ-def word-pred-def word-0-wi word-1-wi*

**lemma** *udvdI*:  
 $0 \leq n \implies \text{uint } b = n * \text{uint } a \implies a \text{ udvd } b$   
 $\langle proof \rangle$

**lemmas** *word-div-no [simp]* =  
*word-div-def [of number-of a number-of b, standard]*

**lemmas** *word-mod-no [simp]* =  
*word-mod-def [of number-of a number-of b, standard]*

**lemmas** *word-less-no [simp]* =  
*word-less-def [of number-of a number-of b, standard]*

**lemmas** *word-le-no [simp]* =  
*word-le-def [of number-of a number-of b, standard]*

**lemmas** *word-sless-no [simp]* =  
*word-sless-def [of number-of a number-of b, standard]*

**lemmas** *word-sle-no [simp]* =  
*word-sle-def [of number-of a number-of b, standard]*

**lemmas** *word-0-wi-Pls* = *word-0-wi [folded Pls-def]*  
**lemmas** *word-0-no* = *word-0-wi-Pls [folded word-no-wi]*

**lemma** *int-one-bin*:  $(1 :: \text{int}) == (\text{Int.Pls BIT bit.B1})$   
 $\langle proof \rangle$

**lemma** *word-1-no*:  
 $(1 :: 'a :: \text{len0 word}) == \text{number-of } (\text{Int.Pls BIT bit.B1})$   
 $\langle proof \rangle$

**lemma** *word-m1-wi*:  $-1 == \text{word-of-int } -1$   
 $\langle proof \rangle$

**lemma** *word-m1-wi-Min*:  $-1 = \text{word-of-int Int.Min}$   
 $\langle proof \rangle$

**lemma** *word-0-bl*: *of-bl*  $\square = 0$   
 $\langle proof \rangle$

**lemma** *word-1-bl*: *of-bl*  $[\text{True}] = 1$



$\langle \text{proof} \rangle$

**lemma** *uint-0* [*simp*] : (*uint* 0 = 0)  
 $\langle \text{proof} \rangle$

**lemma** *of-bl-0* [*simp*] : *of-bl* (*replicate* *n* *False*) = 0  
 $\langle \text{proof} \rangle$

**lemma** *to-bl-0*:  
*to-bl* (0 :: 'a :: len0 word) = *replicate* (*len-of TYPE*('a)) *False*  
 $\langle \text{proof} \rangle$

**lemma** *uint-0-iff*: (*uint* *x* = 0) = (*x* = 0)  
 $\langle \text{proof} \rangle$

**lemma** *unat-0-iff*: (*unat* *x* = 0) = (*x* = 0)  
 $\langle \text{proof} \rangle$

**lemma** *unat-0* [*simp*]: *unat* 0 = 0  
 $\langle \text{proof} \rangle$

**lemma** *size-0-same'*: *size* *w* = 0 ==> *w* = (*v* :: 'a :: len0 word)  
 $\langle \text{proof} \rangle$

**lemmas** *size-0-same* = *size-0-same'* [*folded word-size*]

**lemmas** *unat-eq-0* = *unat-0-iff*  
**lemmas** *unat-eq-zero* = *unat-0-iff*

**lemma** *unat-gt-0*: (0 < *unat* *x*) = (*x* ~ = 0)  
 $\langle \text{proof} \rangle$

**lemma** *ucast-0* [*simp*] : *ucast* 0 = 0  
 $\langle \text{proof} \rangle$

**lemma** *sint-0* [*simp*] : *sint* 0 = 0  
 $\langle \text{proof} \rangle$

**lemma** *scast-0* [*simp*] : *scast* 0 = 0  
 $\langle \text{proof} \rangle$

**lemma** *sint-n1* [*simp*] : *sint* -1 = -1  
 $\langle \text{proof} \rangle$

**lemma** *scast-n1* [*simp*] : *scast* -1 = -1  
 $\langle \text{proof} \rangle$

**lemma** *uint-1* [*simp*] : *uint* (1 :: 'a :: len word) = 1  
 $\langle \text{proof} \rangle$

**lemma** *unat-1* [*simp*] : *unat* (*1* :: '*a* :: *len word*) = *1*  
 ⟨*proof*⟩

**lemma** *ucast-1* [*simp*] : *ucast* (*1* :: '*a* :: *len word*) = *1*  
 ⟨*proof*⟩

**lemmas** *ariths* =  
*bintr-ariths* [*THEN word-ubin.norm-eq-iff* [*THEN iffD1*],  
*folded word-ubin.eq-norm, standard*]

**lemma** *wi-homs*:  
**shows**  
*wi-hom-add*: *word-of-int a* + *word-of-int b* = *word-of-int (a + b)* **and**  
*wi-hom-mult*: *word-of-int a* \* *word-of-int b* = *word-of-int (a \* b)* **and**  
*wi-hom-neg*: − *word-of-int a* = *word-of-int (− a)* **and**  
*wi-hom-succ*: *word-succ (word-of-int a)* = *word-of-int (Int.succ a)* **and**  
*wi-hom-pred*: *word-pred (word-of-int a)* = *word-of-int (Int.pred a)*  
 ⟨*proof*⟩

**lemmas** *wi-hom-syms* = *wi-homs* [*symmetric*]

**lemma** *word-sub-def*: *a* − *b* == *a* + − (*b* :: '*a* :: *len0 word*)  
 ⟨*proof*⟩

**lemmas** *word-diff-minus* = *word-sub-def* [*THEN meta-eq-to-obj-eq, standard*]

**lemma** *word-of-int-sub-hom*:  
 (*word-of-int a*) − *word-of-int b* = *word-of-int (a − b)*  
 ⟨*proof*⟩

**lemmas** *new-word-of-int-homs* =  
*word-of-int-sub-hom wi-homs word-0-wi word-1-wi*

**lemmas** *new-word-of-int-hom-syms* = *new-word-of-int-homs* [*symmetric, standard*]

**lemmas** *word-of-int-hom-syms* =  
*new-word-of-int-hom-syms* [*unfolded succ-def pred-def*]

**lemmas** *word-of-int-homs* =  
*new-word-of-int-homs* [*unfolded succ-def pred-def*]

**lemmas** *word-of-int-add-hom* = *word-of-int-homs* (2)  
**lemmas** *word-of-int-mult-hom* = *word-of-int-homs* (3)  
**lemmas** *word-of-int-minus-hom* = *word-of-int-homs* (4)  
**lemmas** *word-of-int-succ-hom* = *word-of-int-homs* (5)  
**lemmas** *word-of-int-pred-hom* = *word-of-int-homs* (6)

**lemmas** *word-of-int-0-hom* = *word-of-int-homs* (7)

**lemmas** *word-of-int-1-hom* = *word-of-int-homs* (8)

**lemmas** *word-arith-alt* =

*word-sub-wi* [*unfolded succ-def pred-def*, *standard*]

*word-arith-wis* [*unfolded succ-def pred-def*, *standard*]

**lemmas** *word-sub-alt* = *word-arith-alt* (1)

**lemmas** *word-add-alt* = *word-arith-alt* (2)

**lemmas** *word-mult-alt* = *word-arith-alt* (3)

**lemmas** *word-minus-alt* = *word-arith-alt* (4)

**lemmas** *word-succ-alt* = *word-arith-alt* (5)

**lemmas** *word-pred-alt* = *word-arith-alt* (6)

**lemmas** *word-0-alt* = *word-arith-alt* (7)

**lemmas** *word-1-alt* = *word-arith-alt* (8)

### 11.1 Transferring goals from words to ints

**lemma** *word-ths*:

**shows**

*word-succ-p1*: *word-succ a* = *a + 1* **and**

*word-pred-m1*: *word-pred a* = *a - 1* **and**

*word-pred-succ*: *word-pred (word-succ a)* = *a* **and**

*word-succ-pred*: *word-succ (word-pred a)* = *a* **and**

*word-mult-succ*: *word-succ a \* b* = *b + a \* b*

*<proof>*

**lemmas** *uint-cong* = *arg-cong* [**where** *f* = *uint*]

**lemmas** *uint-word-ariths* =

*word-arith-alt* [*THEN trans* [*OF uint-cong int-word-uint*], *standard*]

**lemmas** *uint-word-arith-bintrs* = *uint-word-ariths* [*folded bintrunc-mod2p*]

**lemmas** *sint-word-ariths* = *uint-word-arith-bintrs*

[*THEN uint-sint* [*symmetric*, *THEN trans*],

*unfolded uint-sint bintr-arith1s bintr-ariths*

*len-gt-0* [*THEN bin-sbin-eq-iff*] *word-sbin.norm-Rep*, *standard*]

**lemmas** *uint-div-alt* = *word-div-def*

[*THEN trans* [*OF uint-cong int-word-uint*], *standard*]

**lemmas** *uint-mod-alt* = *word-mod-def*

[*THEN trans* [*OF uint-cong int-word-uint*], *standard*]

**lemma** *word-pred-0-n1*: *word-pred 0* = *word-of-int -1*

*<proof>*

**lemma** *word-pred-0-Min*: *word-pred 0 = word-of-int Int.Min*  
 ⟨*proof*⟩

**lemma** *word-m1-Min*: *− 1 = word-of-int Int.Min*  
 ⟨*proof*⟩

**lemma** *succ-pred-no [simp]*:  
*word-succ (number-of bin) = number-of (Int.succ bin) &*  
*word-pred (number-of bin) = number-of (Int.pred bin)*  
 ⟨*proof*⟩

**lemma** *word-sp-01 [simp]* :  
*word-succ −1 = 0 & word-succ 0 = 1 & word-pred 0 = −1 & word-pred 1 = 0*  
 ⟨*proof*⟩

**lemma** *word-of-int-Ex*:  
 $\exists y. x = \text{word-of-int } y$   
 ⟨*proof*⟩

**lemma** *word-arith-egs*:  
**fixes** *a* :: 'a::len0 word  
**fixes** *b* :: 'a::len0 word  
**shows**  
*word-add-0*:  $0 + a = a$  **and**  
*word-add-0-right*:  $a + 0 = a$  **and**  
*word-mult-1*:  $1 * a = a$  **and**  
*word-mult-1-right*:  $a * 1 = a$  **and**  
*word-add-commute*:  $a + b = b + a$  **and**  
*word-add-assoc*:  $a + b + c = a + (b + c)$  **and**  
*word-add-left-commute*:  $a + (b + c) = b + (a + c)$  **and**  
*word-mult-commute*:  $a * b = b * a$  **and**  
*word-mult-assoc*:  $a * b * c = a * (b * c)$  **and**  
*word-mult-left-commute*:  $a * (b * c) = b * (a * c)$  **and**  
*word-left-distrib*:  $(a + b) * c = a * c + b * c$  **and**  
*word-right-distrib*:  $a * (b + c) = a * b + a * c$  **and**  
*word-left-minus*:  $- a + a = 0$  **and**  
*word-diff-0-right*:  $a - 0 = a$  **and**  
*word-diff-self*:  $a - a = 0$   
 ⟨*proof*⟩

**lemmas** *word-add-ac = word-add-commute word-add-assoc word-add-left-commute*

**lemmas** *word-mult-ac = word-mult-commute word-mult-assoc word-mult-left-commute*

**lemmas** *word-plus-ac0 = word-add-0 word-add-0-right word-add-ac*

**lemmas** *word-times-ac1 = word-mult-1 word-mult-1-right word-mult-ac*

## 11.2 Order on fixed-length words

**lemma** *word-order-trans*:  $x \leq y \implies y \leq z \implies x \leq (z :: 'a :: \text{len0 word})$   
 ⟨proof⟩

**lemma** *word-order-refl*:  $z \leq (z :: 'a :: \text{len0 word})$   
 ⟨proof⟩

**lemma** *word-order-antisym*:  $x \leq y \implies y \leq x \implies x = (y :: 'a :: \text{len0 word})$   
 ⟨proof⟩

**lemma** *word-order-linear*:  
 $y \leq x \mid x \leq (y :: 'a :: \text{len0 word})$   
 ⟨proof⟩

**lemma** *word-zero-le* [simp] :  
 $0 \leq (y :: 'a :: \text{len0 word})$   
 ⟨proof⟩

**instance** *word* :: (len0) semigroup-add  
 ⟨proof⟩

**instance** *word* :: (len0) linorder  
 ⟨proof⟩

**instance** *word* :: (len0) ring  
 ⟨proof⟩

**lemma** *word-m1-ge* [simp] :  $\text{word-pred } 0 \geq y$   
 ⟨proof⟩

**lemmas** *word-n1-ge* [simp] = *word-m1-ge* [simplified *word-sp-01*]

**lemmas** *word-not-simps* [simp] =  
 $\text{word-zero-le } [THEN \text{ leD}] \text{ word-m1-ge } [THEN \text{ leD}] \text{ word-n1-ge } [THEN \text{ leD}]$

**lemma** *word-gt-0*:  $0 < y = (0 \sim (y :: 'a :: \text{len0 word}))$   
 ⟨proof⟩

**lemmas** *word-gt-0-no* [simp] = *word-gt-0* [of *number-of y*, *standard*]

**lemma** *word-sless-alt*:  $(a <_s b) == (\text{sint } a < \text{sint } b)$   
 ⟨proof⟩

**lemma** *word-le-nat-alt*:  $(a \leq b) = (\text{unat } a \leq \text{unat } b)$   
 ⟨proof⟩

**lemma** *word-less-nat-alt*:  $(a < b) = (\text{unat } a < \text{unat } b)$   
 ⟨proof⟩

**lemma** *wi-less*:

$$\begin{aligned} &(\text{word-of-int } n < (\text{word-of-int } m :: 'a :: \text{len0 word})) = \\ & (n \bmod 2 \wedge \text{len-of TYPE}('a) < m \bmod 2 \wedge \text{len-of TYPE}('a)) \\ & \langle \text{proof} \rangle \end{aligned}$$

**lemma** *wi-le*:

$$\begin{aligned} &(\text{word-of-int } n \leq (\text{word-of-int } m :: 'a :: \text{len0 word})) = \\ & (n \bmod 2 \wedge \text{len-of TYPE}('a) \leq m \bmod 2 \wedge \text{len-of TYPE}('a)) \\ & \langle \text{proof} \rangle \end{aligned}$$

**lemma** *udvd-nat-alt*:  $a \text{ udvd } b = (EX n \geq 0. \text{unat } b = n * \text{unat } a)$   
 $\langle \text{proof} \rangle$

**lemma** *udvd-iff-dvd*:  $x \text{ udvd } y \iff \text{unat } x \text{ dvd } \text{unat } y$   
 $\langle \text{proof} \rangle$

**lemmas** *unat-mono* = *word-less-nat-alt* [THEN *iffD1*, *standard*]

**lemma** *word-zero-neq-one*:  $0 < \text{len-of TYPE}('a :: \text{len0}) \implies (0 :: 'a \text{ word}) \sim = 1$   
 $\langle \text{proof} \rangle$

**lemmas** *lenw1-zero-neq-one* = *len-gt-0* [THEN *word-zero-neq-one*]

**lemma** *no-no* [simp]:  $\text{number-of } (\text{number-of } b) = \text{number-of } b$   
 $\langle \text{proof} \rangle$

**lemma** *unat-minus-one*:  $x \sim = 0 \implies \text{unat } (x - 1) = \text{unat } x - 1$   
 $\langle \text{proof} \rangle$

**lemma** *measure-unat*:  $p \sim = 0 \implies \text{unat } (p - 1) < \text{unat } p$   
 $\langle \text{proof} \rangle$

**lemmas** *uint-add-ge0* [simp] =  
*add-nonneg-nonneg* [OF *uint-ge-0 uint-ge-0*, *standard*]

**lemmas** *uint-mult-ge0* [simp] =  
*mult-nonneg-nonneg* [OF *uint-ge-0 uint-ge-0*, *standard*]

**lemma** *uint-sub-lt2p* [simp]:  
 $\text{uint } (x :: 'a :: \text{len0 word}) - \text{uint } (y :: 'b :: \text{len0 word}) <$   
 $2 \wedge \text{len-of TYPE}('a)$   
 $\langle \text{proof} \rangle$

### 11.3 Conditions for the addition (etc) of two words to overflow

**lemma** *uint-add-lem*:

$$\begin{aligned} &(\text{uint } x + \text{uint } y < 2 \wedge \text{len-of TYPE}('a)) = \\ & (\text{uint } (x + y :: 'a :: \text{len0 word}) = \text{uint } x + \text{uint } y) \end{aligned}$$

$\langle \text{proof} \rangle$

**lemma** *uint-mult-lem*:

$(\text{uint } x * \text{uint } y < 2^{\text{len-of TYPE('a)}} =$   
 $(\text{uint } (x * y :: 'a :: \text{len0 word}) = \text{uint } x * \text{uint } y)$   
 $\langle \text{proof} \rangle$

**lemma** *uint-sub-lem*:

$(\text{uint } x \geq \text{uint } y) = (\text{uint } (x - y) = \text{uint } x - \text{uint } y)$   
 $\langle \text{proof} \rangle$

**lemma** *uint-add-le*:  $\text{uint } (x + y) \leq \text{uint } x + \text{uint } y$

$\langle \text{proof} \rangle$

**lemma** *uint-sub-ge*:  $\text{uint } (x - y) \geq \text{uint } x - \text{uint } y$

$\langle \text{proof} \rangle$

**lemmas** *uint-sub-if'* =

*trans* [*OF uint-word-ariths*(1) *mod-sub-if-z, simplified, standard*]

**lemmas** *uint-plus-if'* =

*trans* [*OF uint-word-ariths*(2) *mod-add-if-z, simplified, standard*]

## 11.4 Definition of uint\_arith

**lemma** *word-of-int-inverse*:

$\text{word-of-int } r = a \implies 0 \leq r \implies r < 2^{\text{len-of TYPE('a)}} \implies$   
 $\text{uint } (a :: 'a :: \text{len0 word}) = r$   
 $\langle \text{proof} \rangle$

**lemma** *uint-split*:

**fixes**  $x :: 'a :: \text{len0 word}$

**shows**  $P (\text{uint } x) =$

$(\text{ALL } i. \text{word-of-int } i = x \ \& \ 0 \leq i \ \& \ i < 2^{\text{len-of TYPE('a)}} \implies P \ i)$

$\langle \text{proof} \rangle$

**lemma** *uint-split-asm*:

**fixes**  $x :: 'a :: \text{len0 word}$

**shows**  $P (\text{uint } x) =$

$(\sim (\text{EX } i. \text{word-of-int } i = x \ \& \ 0 \leq i \ \& \ i < 2^{\text{len-of TYPE('a)}} \ \& \ \sim P \ i))$

$\langle \text{proof} \rangle$

**lemmas** *uint-splits* = *uint-split uint-split-asm*

**lemmas** *uint-arith-simps* =

*word-le-def word-less-alt*

*word-uint.Rep-inject* [*symmetric*]

*uint-sub-if' uint-plus-if'*

**lemma** *power-False-cong*:  $\text{False} ==> a \wedge b = c \wedge d$   
 ⟨proof⟩

⟨ML⟩

## 11.5 More on overflows and monotonicity

**lemma** *no-plus-overflow-uint-size*:  
 $((x :: 'a :: \text{len0 word}) <= x + y) = (\text{uint } x + \text{uint } y < 2 \wedge \text{size } x)$   
 ⟨proof⟩

**lemmas** *no-olen-add* = *no-plus-overflow-uint-size* [unfolded word-size]

**lemma** *no-ulen-sub*:  $((x :: 'a :: \text{len0 word}) >= x - y) = (\text{uint } y <= \text{uint } x)$   
 ⟨proof⟩

**lemma** *no-olen-add'*:  
 fixes  $x :: 'a :: \text{len0 word}$   
 shows  $(x \leq y + x) = (\text{uint } y + \text{uint } x < 2 \wedge \text{len-of TYPE('a)})$   
 ⟨proof⟩

**lemmas** *olen-add-eqv* = *trans* [OF *no-olen-add no-olen-add'* [symmetric], standard]

**lemmas** *uint-plus-simple-iff* = *trans* [OF *no-olen-add uint-add-lem*, standard]  
**lemmas** *uint-plus-simple* = *uint-plus-simple-iff* [THEN *iffD1*, standard]  
**lemmas** *uint-minus-simple-iff* = *trans* [OF *no-ulen-sub uint-sub-lem*, standard]  
**lemmas** *uint-minus-simple-alt* = *uint-sub-lem* [folded word-le-def]  
**lemmas** *word-sub-le-iff* = *no-ulen-sub* [folded word-le-def]  
**lemmas** *word-sub-le* = *word-sub-le-iff* [THEN *iffD2*, standard]

**lemma** *word-less-sub1*:  
 $(x :: 'a :: \text{len word}) \sim= 0 ==> (1 < x) = (0 < x - 1)$   
 ⟨proof⟩

**lemma** *word-le-sub1*:  
 $(x :: 'a :: \text{len word}) \sim= 0 ==> (1 <= x) = (0 <= x - 1)$   
 ⟨proof⟩

**lemma** *sub-wrap-lt*:  
 $((x :: 'a :: \text{len0 word}) < x - z) = (x < z)$   
 ⟨proof⟩

**lemma** *sub-wrap*:  
 $((x :: 'a :: \text{len0 word}) <= x - z) = (z = 0 \mid x < z)$   
 ⟨proof⟩

**lemma** *plus-minus-not-NULL-ab*:



$(x :: 'a :: \text{len0 word}) <= ab - c ==> c <= ab ==> c \sim 0 ==> x + c \sim 0$   
 $\langle \text{proof} \rangle$

**lemma** *plus-minus-no-overflow-ab*:

$(x :: 'a :: \text{len0 word}) <= ab - c ==> c <= ab ==> x <= x + c$   
 $\langle \text{proof} \rangle$

**lemma** *le-minus'*:

$(a :: 'a :: \text{len0 word}) + c <= b ==> a <= a + c ==> c <= b - a$   
 $\langle \text{proof} \rangle$

**lemma** *le-plus'*:

$(a :: 'a :: \text{len0 word}) <= b ==> c <= b - a ==> a + c <= b$   
 $\langle \text{proof} \rangle$

**lemmas** *le-plus = le-plus'* [rotated]

**lemmas** *le-minus = leD* [THEN *thin-rl*, THEN *le-minus'*, standard]

**lemma** *word-plus-mono-right*:

$(y :: 'a :: \text{len0 word}) <= z ==> x <= x + z ==> x + y <= x + z$   
 $\langle \text{proof} \rangle$

**lemma** *word-less-minus-cancel*:

$y - x < z - x ==> x <= z ==> (y :: 'a :: \text{len0 word}) < z$   
 $\langle \text{proof} \rangle$

**lemma** *word-less-minus-mono-left*:

$(y :: 'a :: \text{len0 word}) < z ==> x <= y ==> y - x < z - x$   
 $\langle \text{proof} \rangle$

**lemma** *word-less-minus-mono*:

$a < c ==> d < b ==> a - b < a ==> c - d < c$   
 $==> a - b < c - (d :: 'a :: \text{len word})$   
 $\langle \text{proof} \rangle$

**lemma** *word-le-minus-cancel*:

$y - x <= z - x ==> x <= z ==> (y :: 'a :: \text{len0 word}) <= z$   
 $\langle \text{proof} \rangle$

**lemma** *word-le-minus-mono-left*:

$(y :: 'a :: \text{len0 word}) <= z ==> x <= y ==> y - x <= z - x$   
 $\langle \text{proof} \rangle$

**lemma** *word-le-minus-mono*:

$a <= c ==> d <= b ==> a - b <= a ==> c - d <= c$   
 $==> a - b <= c - (d :: 'a :: \text{len word})$   
 $\langle \text{proof} \rangle$

**lemma** *plus-le-left-cancel-wrap*:

$$(x :: 'a :: \text{len0 word}) + y' < x \implies x + y < x \implies (x + y' < x + y) = (y' < y)$$

*<proof>*

**lemma** *plus-le-left-cancel-nowrap*:

$$(x :: 'a :: \text{len0 word}) <= x + y' \implies x <= x + y \implies (x + y' < x + y) = (y' < y)$$

*<proof>*

**lemma** *word-plus-mono-right2*:

$$(a :: 'a :: \text{len0 word}) <= a + b \implies c <= b \implies a <= a + c$$

*<proof>*

**lemma** *word-less-add-right*:

$$(x :: 'a :: \text{len0 word}) < y - z \implies z <= y \implies x + z < y$$

*<proof>*

**lemma** *word-less-sub-right*:

$$(x :: 'a :: \text{len0 word}) < y + z \implies y <= x \implies x - y < z$$

*<proof>*

**lemma** *word-le-plus-either*:

$$(x :: 'a :: \text{len0 word}) <= y \mid x <= z \implies y <= y + z \implies x <= y + z$$

*<proof>*

**lemma** *word-less-nowrapI*:

$$(x :: 'a :: \text{len0 word}) < z - k \implies k <= z \implies 0 < k \implies x < x + k$$

*<proof>*

**lemma** *inc-le*:  $(i :: 'a :: \text{len word}) < m \implies i + 1 <= m$

*<proof>*

**lemma** *inc-i*:

$$(1 :: 'a :: \text{len word}) <= i \implies i < m \implies 1 <= (i + 1) \ \& \ i + 1 <= m$$

*<proof>*

**lemma** *udvd-incr-lem*:

$$\begin{aligned} up < uq &\implies up = ua + n * \text{uint } K \implies \\ uq = ua + n' * \text{uint } K &\implies up + \text{uint } K <= uq \end{aligned}$$

*<proof>*

**lemma** *udvd-incr'*:

$$\begin{aligned} p < q &\implies \text{uint } p = ua + n * \text{uint } K \implies \\ \text{uint } q = ua + n' * \text{uint } K &\implies p + K <= q \end{aligned}$$

*<proof>*

**lemma** *udvd-decr'*:

$$p < q \implies \text{uint } p = ua + n * \text{uint } K \implies$$

$uint\ q = ua + n' * uint\ K ==> p <= q - K$   
 $\langle proof \rangle$

**lemmas**  $udvd-incr-lem0 = udvd-incr-lem$  [where  $ua=0$ , simplified]  
**lemmas**  $udvd-incr0 = udvd-incr'$  [where  $ua=0$ , simplified]  
**lemmas**  $udvd-decr0 = udvd-decr'$  [where  $ua=0$ , simplified]

**lemma**  $udvd-minus-le'$ :  
 $xy < k ==> z\ udvd\ xy ==> z\ udvd\ k ==> xy <= k - z$   
 $\langle proof \rangle$

**lemma**  $udvd-incr2-K$ :  
 $p < a + s ==> a <= a + s ==> K\ udvd\ s ==> K\ udvd\ p - a ==> a <=$   
 $p ==>$   
 $0 < K ==> p <= p + K \ \& \ p + K <= a + s$   
 $\langle proof \rangle$

**lemma**  $word-succ-rbl$ :  
 $to-bl\ w = bl ==> to-bl\ (word-succ\ w) = (rev\ (rbl-succ\ (rev\ bl)))$   
 $\langle proof \rangle$

**lemma**  $word-pred-rbl$ :  
 $to-bl\ w = bl ==> to-bl\ (word-pred\ w) = (rev\ (rbl-pred\ (rev\ bl)))$   
 $\langle proof \rangle$

**lemma**  $word-add-rbl$ :  
 $to-bl\ v = vbl ==> to-bl\ w = wbl ==>$   
 $to-bl\ (v + w) = (rev\ (rbl-add\ (rev\ vbl)\ (rev\ wbl)))$   
 $\langle proof \rangle$

**lemma**  $word-mult-rbl$ :  
 $to-bl\ v = vbl ==> to-bl\ w = wbl ==>$   
 $to-bl\ (v * w) = (rev\ (rbl-mult\ (rev\ vbl)\ (rev\ wbl)))$   
 $\langle proof \rangle$

**lemma**  $rtb-rbl-ariths$ :  
 $rev\ (to-bl\ w) = ys \implies rev\ (to-bl\ (word-succ\ w)) = rbl-succ\ ys$

$rev\ (to-bl\ w) = ys \implies rev\ (to-bl\ (word-pred\ w)) = rbl-pred\ ys$

$[| rev\ (to-bl\ v) = ys; rev\ (to-bl\ w) = xs |]$   
 $==> rev\ (to-bl\ (v * w)) = rbl-mult\ ys\ xs$

$[| rev\ (to-bl\ v) = ys; rev\ (to-bl\ w) = xs |]$   
 $==> rev\ (to-bl\ (v + w)) = rbl-add\ ys\ xs$   
 $\langle proof \rangle$

## 11.6 Arithmetic type class instantiations

**instance** *word* :: (*len0*) *comm-monoid-add*  $\langle$ *proof* $\rangle$

**instance** *word* :: (*len0*) *comm-monoid-mult*  
 $\langle$ *proof* $\rangle$

**instance** *word* :: (*len0*) *comm-semiring*  
 $\langle$ *proof* $\rangle$

**instance** *word* :: (*len0*) *ab-group-add*  $\langle$ *proof* $\rangle$

**instance** *word* :: (*len0*) *comm-ring*  $\langle$ *proof* $\rangle$

**instance** *word* :: (*len*) *comm-semiring-1*  
 $\langle$ *proof* $\rangle$

**instance** *word* :: (*len*) *comm-ring-1*  $\langle$ *proof* $\rangle$

**instance** *word* :: (*len0*) *comm-semiring-0*  $\langle$ *proof* $\rangle$

**instance** *word* :: (*len0*) *order*  $\langle$ *proof* $\rangle$

**instance** *word* :: (*len*) *recpower*  
 $\langle$ *proof* $\rangle$

**lemma** *zero-bintrunc*:  
 $\text{iszero } (\text{number-of } x :: 'a :: \text{len } \text{word}) =$   
 $(\text{bintrunc } (\text{len-of TYPE('a)}) x = \text{Int.Pl})$   
 $\langle$ *proof* $\rangle$

**lemmas** *word-le-0-iff* [*simp*] =  
 $\text{word-zero-le } [\text{THEN } \text{leD}, \text{ THEN } \text{linorder-antisym-conv1}]$

**lemma** *word-of-nat*:  $\text{of-nat } n = \text{word-of-int } (\text{int } n)$   
 $\langle$ *proof* $\rangle$

**lemma** *word-of-int*:  $\text{of-int} = \text{word-of-int}$   
 $\langle$ *proof* $\rangle$

**lemma** *word-of-int-nat*:  
 $0 \leq x \implies \text{word-of-int } x = \text{of-nat } (\text{nat } x)$   
 $\langle$ *proof* $\rangle$

**lemma** *word-number-of-eq*:  
 $\text{number-of } w = (\text{of-int } w :: 'a :: \text{len } \text{word})$   
 $\langle$ *proof* $\rangle$

**instance** *word* :: (*len*) *number-ring*

*<proof>*

**lemma** *iszero-word-no* [simp] :  
 $\text{iszero } (\text{number-of bin} :: 'a :: \text{len word}) =$   
 $\text{iszero } (\text{number-of } (\text{bintrunc } (\text{len-of TYPE('a)}) \text{ bin}) :: \text{int})$   
*<proof>*

## 11.7 Word and nat

**lemma** *td-ext-unat'*:  
 $n = \text{len-of TYPE } ('a :: \text{len}) ==>$   
 $\text{td-ext } (\text{unat} :: 'a \text{ word} ==> \text{nat}) \text{ of-nat}$   
 $(\text{unats } n) (\%i. i \bmod 2^n)$   
*<proof>*

**lemmas** *td-ext-unat* = refl [THEN *td-ext-unat'*]  
**lemmas** *unat-of-nat* = *td-ext-unat* [THEN *td-ext.eq-norm*, *standard*]

**interpretation** *word-unat*:  
 $\text{td-ext } [\text{unat} :: 'a :: \text{len word} ==> \text{nat}]$   
 $\text{of-nat}$   
 $\text{unats } (\text{len-of TYPE } ('a :: \text{len}))$   
 $\%i. i \bmod 2^{\text{len-of TYPE } ('a :: \text{len})}$   
*<proof>*

**lemmas** *td-unat* = *word-unat.td-thm*

**lemmas** *unat-lt2p* [iff] = *word-unat.Rep* [unfolded *unats-def mem-Collect-eq*]

**lemma** *unat-le*:  $y \leq \text{unat } (z :: 'a :: \text{len word}) ==> y : \text{unats } (\text{len-of TYPE } ('a))$   
*<proof>*

**lemma** *word-nchotomy*:  
 $\text{ALL } w. \text{EX } n. (w :: 'a :: \text{len word}) = \text{of-nat } n \ \& \ n < 2^{\text{len-of TYPE } ('a)}$   
*<proof>*

**lemma** *of-nat-eq*:  
**fixes**  $w :: 'a :: \text{len word}$   
**shows**  $(\text{of-nat } n = w) = (\exists q. n = \text{unat } w + q * 2^{\text{len-of TYPE } ('a)})$   
*<proof>*

**lemma** *of-nat-eq-size*:  
 $(\text{of-nat } n = w) = (\text{EX } q. n = \text{unat } w + q * 2^{\text{size } w})$   
*<proof>*

**lemma** *of-nat-0*:  
 $(\text{of-nat } m = (0 :: 'a :: \text{len word})) = (\exists q. m = q * 2^{\text{len-of TYPE } ('a)})$   
*<proof>*

**lemmas** *of-nat-2p* = *mult-1* [*symmetric*, *THEN iffD2* [*OF of-nat-0 exI*]]

**lemma** *of-nat-gt-0*: *of-nat k*  $\sim$  *0*  $\implies$  *0* < *k*  
 ⟨*proof*⟩

**lemma** *of-nat-neq-0*:  
*0* < *k*  $\implies$  *k* < *2*  $\wedge$  *len-of TYPE* (*'a* :: *len*)  $\implies$  *of-nat k*  $\sim$  (*0* :: *'a word*)  
 ⟨*proof*⟩

**lemma** *Abs-fnat-hom-add*:  
*of-nat a* + *of-nat b* = *of-nat (a + b)*  
 ⟨*proof*⟩

**lemma** *Abs-fnat-hom-mult*:  
*of-nat a* \* *of-nat b* = (*of-nat (a \* b)* :: *'a* :: *len word*)  
 ⟨*proof*⟩

**lemma** *Abs-fnat-hom-Suc*:  
*word-succ (of-nat a)* = *of-nat (Suc a)*  
 ⟨*proof*⟩

**lemma** *Abs-fnat-hom-0*: (*0*::*'a*::*len word*) = *of-nat 0*  
 ⟨*proof*⟩

**lemma** *Abs-fnat-hom-1*: (*1*::*'a*::*len word*) = *of-nat (Suc 0)*  
 ⟨*proof*⟩

**lemmas** *Abs-fnat-homs* =  
*Abs-fnat-hom-add* *Abs-fnat-hom-mult* *Abs-fnat-hom-Suc*  
*Abs-fnat-hom-0* *Abs-fnat-hom-1*

**lemma** *word-arith-nat-add*:  
*a* + *b* = *of-nat (unat a + unat b)*  
 ⟨*proof*⟩

**lemma** *word-arith-nat-mult*:  
*a* \* *b* = *of-nat (unat a \* unat b)*  
 ⟨*proof*⟩

**lemma** *word-arith-nat-Suc*:  
*word-succ a* = *of-nat (Suc (unat a))*  
 ⟨*proof*⟩

**lemma** *word-arith-nat-div*:  
*a div b* = *of-nat (unat a div unat b)*  
 ⟨*proof*⟩

**lemma** *word-arith-nat-mod*:

$a \bmod b = \text{of-nat } (\text{unat } a \bmod \text{unat } b)$   
 $\langle \text{proof} \rangle$

**lemmas** *word-arith-nat-defs* =  
*word-arith-nat-add word-arith-nat-mult*  
*word-arith-nat-Suc Abs-fnat-hom-0*  
*Abs-fnat-hom-1 word-arith-nat-div*  
*word-arith-nat-mod*

**lemmas** *unat-cong* = *arg-cong* [where  $f = \text{unat}$ ]

**lemmas** *unat-word-ariths* = *word-arith-nat-defs*  
[*THEN trans* [*OF unat-cong unat-of-nat*], *standard*]

**lemmas** *word-sub-less-iff* = *word-sub-le-iff*  
[*simplified linorder-not-less* [*symmetric*], *simplified*]

**lemma** *unat-add-lem*:  
 $(\text{unat } x + \text{unat } y < 2 \wedge \text{len-of TYPE('a)}) =$   
 $(\text{unat } (x + y :: 'a :: \text{len word}) = \text{unat } x + \text{unat } y)$   
 $\langle \text{proof} \rangle$

**lemma** *unat-mult-lem*:  
 $(\text{unat } x * \text{unat } y < 2 \wedge \text{len-of TYPE('a)}) =$   
 $(\text{unat } (x * y :: 'a :: \text{len word}) = \text{unat } x * \text{unat } y)$   
 $\langle \text{proof} \rangle$

**lemmas** *unat-plus-if'* =  
*trans* [*OF unat-word-ariths(1) mod-nat-add, simplified, standard*]

**lemma** *le-no-overflow*:  
 $x \leq b \implies a \leq a + b \implies x \leq a + (b :: 'a :: \text{len0 word})$   
 $\langle \text{proof} \rangle$

**lemmas** *un-ui-le* = *trans*  
[*OF word-le-nat-alt* [*symmetric*]  
*word-le-def*,  
*standard*]

**lemma** *unat-sub-if-size*:  
 $\text{unat } (x - y) = (\text{if } \text{unat } y \leq \text{unat } x$   
 $\text{then } \text{unat } x - \text{unat } y$   
 $\text{else } \text{unat } x + 2 \wedge \text{size } x - \text{unat } y)$   
 $\langle \text{proof} \rangle$

**lemmas** *unat-sub-if'* = *unat-sub-if-size* [*unfolded word-size*]

**lemma** *unat-div*:  $\text{unat } ((x :: 'a :: \text{len word}) \text{ div } y) = \text{unat } x \text{ div } \text{unat } y$   
 $\langle \text{proof} \rangle$

**lemma** *unat-mod*:  $\text{unat } ((x :: 'a :: \text{len word}) \text{ mod } y) = \text{unat } x \text{ mod } \text{unat } y$   
*<proof>*

**lemma** *uint-div*:  $\text{uint } ((x :: 'a :: \text{len word}) \text{ div } y) = \text{uint } x \text{ div } \text{uint } y$   
*<proof>*

**lemma** *uint-mod*:  $\text{uint } ((x :: 'a :: \text{len word}) \text{ mod } y) = \text{uint } x \text{ mod } \text{uint } y$   
*<proof>*

## 11.8 Definition of unat\_arith tactic

**lemma** *unat-split*:  
**fixes**  $x :: 'a :: \text{len word}$   
**shows**  $P (\text{unat } x) =$   
 $(\text{ALL } n. \text{ of-nat } n = x \ \& \ n < 2^{\text{len-of TYPE('a)}} \longrightarrow P \ n)$   
*<proof>*

**lemma** *unat-split-asm*:  
**fixes**  $x :: 'a :: \text{len word}$   
**shows**  $P (\text{unat } x) =$   
 $(\sim (\text{EX } n. \text{ of-nat } n = x \ \& \ n < 2^{\text{len-of TYPE('a)}} \ \& \ \sim P \ n))$   
*<proof>*

**lemmas** *of-nat-inverse* =  
 $\text{word-unat.Abs-inverse'}$  [*rotated, unfolded unats-def, simplified*]

**lemmas** *unat-splits* = *unat-split unat-split-asm*

**lemmas** *unat-arith-simps* =  
 $\text{word-le-nat-alt}$   $\text{word-less-nat-alt}$   
 $\text{word-unat.Rep-inject}$  [*symmetric*]  
 $\text{unat-sub-if'}$   $\text{unat-plus-if'}$   $\text{unat-div}$   $\text{unat-mod}$

*<ML>*

**lemma** *no-plus-overflow-unat-size*:  
 $((x :: 'a :: \text{len word}) \leq x + y) = (\text{unat } x + \text{unat } y < 2^{\text{size } x})$   
*<proof>*

**lemma** *unat-sub*:  $b \leq a \implies \text{unat } (a - b) = \text{unat } a - \text{unat } (b :: 'a :: \text{len word})$   
*<proof>*

**lemmas** *no-olen-add-nat* = *no-plus-overflow-unat-size* [*unfolded word-size*]

**lemmas** *unat-plus-simple* = *trans* [*OF no-olen-add-nat unat-add-lem, standard*]

**lemma** *word-div-mult*:



$(0 :: 'a :: \text{len word}) < y ==> \text{unat } x * \text{unat } y < 2 \wedge \text{len-of TYPE}('a) ==>$   
 $x * y \text{ div } y = x$   
 $\langle \text{proof} \rangle$

**lemma** *div-lt'*:  $(i :: 'a :: \text{len word}) <= k \text{ div } x ==>$   
 $\text{unat } i * \text{unat } x < 2 \wedge \text{len-of TYPE}('a)$   
 $\langle \text{proof} \rangle$

**lemmas** *div-lt''* = *order-less-imp-le* [THEN *div-lt'*]

**lemma** *div-lt-mult*:  $(i :: 'a :: \text{len word}) < k \text{ div } x ==> 0 < x ==> i * x < k$   
 $\langle \text{proof} \rangle$

**lemma** *div-le-mult*:  
 $(i :: 'a :: \text{len word}) <= k \text{ div } x ==> 0 < x ==> i * x <= k$   
 $\langle \text{proof} \rangle$

**lemma** *div-lt-uint'*:  
 $(i :: 'a :: \text{len word}) <= k \text{ div } x ==> \text{uint } i * \text{uint } x < 2 \wedge \text{len-of TYPE}('a)$   
 $\langle \text{proof} \rangle$

**lemmas** *div-lt-uint''* = *order-less-imp-le* [THEN *div-lt-uint'*]

**lemma** *word-le-exists'*:  
 $(x :: 'a :: \text{len0 word}) <= y ==>$   
 $(\text{EX } z. y = x + z \ \& \ \text{uint } x + \text{uint } z < 2 \wedge \text{len-of TYPE}('a))$   
 $\langle \text{proof} \rangle$

**lemmas** *plus-minus-not-NULL* = *order-less-imp-le* [THEN *plus-minus-not-NULL-ab*]

**lemmas** *plus-minus-no-overflow* =  
*order-less-imp-le* [THEN *plus-minus-no-overflow-ab*]

**lemmas** *mcs* = *word-less-minus-cancel word-less-minus-mono-left*  
*word-le-minus-cancel word-le-minus-mono-left*

**lemmas** *word-l-diffs* = *mcs* [where  $y = w + x$ , *unfolded add-diff-cancel*, *standard*]  
**lemmas** *word-diff-ls* = *mcs* [where  $z = w + x$ , *unfolded add-diff-cancel*, *standard*]  
**lemmas** *word-plus-mcs* = *word-diff-ls*  
[where  $y = v + x$ , *unfolded add-diff-cancel*, *standard*]

**lemmas** *le-unat-woi* = *unat-le* [THEN *word-unat.Abs-inverse*]

**lemmas** *thd* = *refl* [THEN [2] *split-div-lemma* [THEN *iffD2*], THEN *conjunct1*]

**lemma** *thd1*:  
 $a \text{ div } b * b \leq (a :: \text{nat})$   
 $\langle \text{proof} \rangle$

**lemmas** *uno-simps* [*THEN le-unat-voi, standard*] =  
*mod-le-divisor div-le-dividend thd1*

**lemma** *word-mod-div-equality*:  
 $(n \text{ div } b) * b + (n \text{ mod } b) = (n :: 'a :: \text{len word})$   
*<proof>*

**lemma** *word-div-mult-le*:  $a \text{ div } b * b \leq (a :: 'a :: \text{len word})$   
*<proof>*

**lemma** *word-mod-less-divisor*:  $0 < n \implies m \text{ mod } n < (n :: 'a :: \text{len word})$   
*<proof>*

**lemma** *word-of-int-power-hom*:  
 $\text{word-of-int } a ^ n = (\text{word-of-int } (a ^ n) :: 'a :: \text{len word})$   
*<proof>*

**lemma** *word-arith-power-alt*:  
 $a ^ n = (\text{word-of-int } (\text{uint } a ^ n) :: 'a :: \text{len word})$   
*<proof>*

**lemma** *of-bl-length-less*:  
 $\text{length } x = k \implies k < \text{len-of TYPE('a)} \implies (\text{of-bl } x :: 'a :: \text{len word}) < 2 ^ k$   
*<proof>*

## 11.9 Cardinality, finiteness of set of words

**lemmas** *card-lessThan'* = *card-lessThan* [*unfolded lessThan-def*]

**lemmas** *card-eq* = *word-unat.Abs-inj-on* [*THEN card-image,*  
*unfolded word-unat.image, unfolded unats-def, standard*]

**lemmas** *card-word* = *trans* [*OF card-eq card-lessThan', standard*]

**lemma** *finite-word-UNIV*: *finite* (*UNIV* :: *'a* :: *len word set*)  
*<proof>*

**lemma** *card-word-size*:  
 $\text{card } (\text{UNIV} :: 'a :: \text{len word set}) = (2 ^ \text{size } (x :: 'a \text{ word}))$   
*<proof>*

**end**

## 12 WordBitwise: Bitwise Operations on Words

**theory** *WordBitwise*  
**imports** *WordArith*

**begin**

**lemmas** *bin-log-bintrs* = *bin-trunc-not bin-trunc-xor bin-trunc-and bin-trunc-or*

**lemmas** *wils1* = *bin-log-bintrs* [THEN *word-ubin.norm-eq-iff* [THEN *iffD1*],  
folded *word-ubin.eq-norm*, THEN *eq-reflection*, standard]

**lemmas** *word-log-binary-defs* =  
*word-and-def word-or-def word-xor-def*

**lemmas** *word-no-log-defs* [*simp*] =  
*word-not-def* [where *a=number-of a*,  
unfolded *word-no-wi wils1*, folded *word-no-wi*, standard]  
*word-log-binary-defs* [where *a=number-of a* and *b=number-of b*,  
unfolded *word-no-wi wils1*, folded *word-no-wi*, standard]

**lemmas** *word-wi-log-defs* = *word-no-log-defs* [unfolded *word-no-wi*]

**lemma** *uint-or*: *uint (x OR y) = (uint x) OR (uint y)*  
⟨*proof*⟩

**lemma** *uint-and*: *uint (x AND y) = (uint x) AND (uint y)*  
⟨*proof*⟩

**lemma** *word-ops-nth-size*:  
*n < size (x::'a::len0 word) ==>*  
*(x OR y) !! n = (x !! n | y !! n) &*  
*(x AND y) !! n = (x !! n & y !! n) &*  
*(x XOR y) !! n = (x !! n ~ y !! n) &*  
*(NOT x) !! n = (~ x !! n)*  
⟨*proof*⟩

**lemma** *word-ao-nth*:  
**fixes** *x :: 'a::len0 word*  
**shows** *(x OR y) !! n = (x !! n | y !! n) &*  
*(x AND y) !! n = (x !! n & y !! n)*  
⟨*proof*⟩

**lemmas** *bwsimps* =  
*word-of-int-homs(2)*  
*word-0-wi-Pls*  
*word-m1-wi-Min*  
*word-wi-log-defs*

**lemma** *word-bw-assocs*:

**fixes**  $x :: 'a::len0$  *word*

**shows**

$(x \text{ AND } y) \text{ AND } z = x \text{ AND } y \text{ AND } z$

$(x \text{ OR } y) \text{ OR } z = x \text{ OR } y \text{ OR } z$

$(x \text{ XOR } y) \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$

$\langle \text{proof} \rangle$

**lemma** *word-bw-comms*:

**fixes**  $x :: 'a::len0$  *word*

**shows**

$x \text{ AND } y = y \text{ AND } x$

$x \text{ OR } y = y \text{ OR } x$

$x \text{ XOR } y = y \text{ XOR } x$

$\langle \text{proof} \rangle$

**lemma** *word-bw-lcs*:

**fixes**  $x :: 'a::len0$  *word*

**shows**

$y \text{ AND } x \text{ AND } z = x \text{ AND } y \text{ AND } z$

$y \text{ OR } x \text{ OR } z = x \text{ OR } y \text{ OR } z$

$y \text{ XOR } x \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$

$\langle \text{proof} \rangle$

**lemma** *word-log-esimps* [*simp*]:

**fixes**  $x :: 'a::len0$  *word*

**shows**

$x \text{ AND } 0 = 0$

$x \text{ AND } -1 = x$

$x \text{ OR } 0 = x$

$x \text{ OR } -1 = -1$

$x \text{ XOR } 0 = x$

$x \text{ XOR } -1 = \text{NOT } x$

$0 \text{ AND } x = 0$

$-1 \text{ AND } x = x$

$0 \text{ OR } x = x$

$-1 \text{ OR } x = -1$

$0 \text{ XOR } x = x$

$-1 \text{ XOR } x = \text{NOT } x$

$\langle \text{proof} \rangle$

**lemma** *word-not-dist*:

**fixes**  $x :: 'a::len0$  *word*

**shows**

$\text{NOT } (x \text{ OR } y) = \text{NOT } x \text{ AND } \text{NOT } y$

$\text{NOT } (x \text{ AND } y) = \text{NOT } x \text{ OR } \text{NOT } y$

$\langle \text{proof} \rangle$

**lemma** *word-bw-same*:

**fixes**  $x :: 'a::len0 \text{ word}$

**shows**

$x \text{ AND } x = x$

$x \text{ OR } x = x$

$x \text{ XOR } x = 0$

$\langle \text{proof} \rangle$

**lemma** *word-ao-absorbs* [simp]:

**fixes**  $x :: 'a::len0 \text{ word}$

**shows**

$x \text{ AND } (y \text{ OR } x) = x$

$x \text{ OR } y \text{ AND } x = x$

$x \text{ AND } (x \text{ OR } y) = x$

$y \text{ AND } x \text{ OR } x = x$

$(y \text{ OR } x) \text{ AND } x = x$

$x \text{ OR } x \text{ AND } y = x$

$(x \text{ OR } y) \text{ AND } x = x$

$x \text{ AND } y \text{ OR } x = x$

$\langle \text{proof} \rangle$

**lemma** *word-not-not* [simp]:

$\text{NOT NOT } (x :: 'a::len0 \text{ word}) = x$

$\langle \text{proof} \rangle$

**lemma** *word-ao-dist*:

**fixes**  $x :: 'a::len0 \text{ word}$

**shows**  $(x \text{ OR } y) \text{ AND } z = x \text{ AND } z \text{ OR } y \text{ AND } z$

$\langle \text{proof} \rangle$

**lemma** *word-oa-dist*:

**fixes**  $x :: 'a::len0 \text{ word}$

**shows**  $x \text{ AND } y \text{ OR } z = (x \text{ OR } z) \text{ AND } (y \text{ OR } z)$

$\langle \text{proof} \rangle$

**lemma** *word-add-not* [simp]:

**fixes**  $x :: 'a::len0 \text{ word}$

**shows**  $x + \text{NOT } x = -1$

$\langle \text{proof} \rangle$

**lemma** *word-plus-and-or* [simp]:

**fixes**  $x :: 'a::len0 \text{ word}$

**shows**  $(x \text{ AND } y) + (x \text{ OR } y) = x + y$

$\langle \text{proof} \rangle$

**lemma** *leoa*:

**fixes**  $x :: 'a::len0 \text{ word}$

**shows**  $(w = (x \text{ OR } y)) ==> (y = (w \text{ AND } y)) \langle \text{proof} \rangle$

**lemma** *leao*:

**fixes**  $x' :: 'a::len0 \text{ word}$   
**shows**  $(w' = (x' \text{ AND } y')) ==> (x' = (x' \text{ OR } w')) \langle \text{proof} \rangle$

**lemmas**  $\text{word-ao-equiv} = \text{leao} [\text{COMP } \text{leoa} [\text{COMP } \text{iffI}]]$

**lemma**  $\text{le-word-or2}: x \leq x \text{ OR } (y :: 'a::len0 \text{ word})$   
 $\langle \text{proof} \rangle$

**lemmas**  $\text{le-word-or1} = \text{xtr3} [\text{OF } \text{word-bw-comms } (2) \text{ le-word-or2, standard}]$

**lemmas**  $\text{word-and-le1} =$   
 $\text{xtr3} [\text{OF } \text{word-ao-absorbs } (4) [\text{symmetric}] \text{ le-word-or2, standard}]$

**lemmas**  $\text{word-and-le2} =$   
 $\text{xtr3} [\text{OF } \text{word-ao-absorbs } (8) [\text{symmetric}] \text{ le-word-or2, standard}]$

**lemma**  $\text{bl-word-not}: \text{to-bl } (\text{NOT } w) = \text{map Not } (\text{to-bl } w)$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bl-word-xor}: \text{to-bl } (v \text{ XOR } w) = \text{map2 op } \sim = (\text{to-bl } v) (\text{to-bl } w)$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bl-word-or}: \text{to-bl } (v \text{ OR } w) = \text{map2 op } | (\text{to-bl } v) (\text{to-bl } w)$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{bl-word-and}: \text{to-bl } (v \text{ AND } w) = \text{map2 op } \& (\text{to-bl } v) (\text{to-bl } w)$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{word-lsb-alt}: \text{lsb } (w :: 'a::len0 \text{ word}) = \text{test-bit } w \ 0$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{word-lsb-1-0}: \text{lsb } (1 :: 'a::len \text{ word}) \& \sim \text{lsb } (0 :: 'b::len0 \text{ word})$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{word-lsb-last}: \text{lsb } (w :: 'a::len \text{ word}) = \text{last } (\text{to-bl } w)$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{word-lsb-int}: \text{lsb } w = (\text{uint } w \bmod 2 = 1)$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{word-msb-sint}: \text{msb } w = (\text{sint } w < 0)$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{word-msb-no}':$   
 $w = \text{number-of bin} ==> \text{msb } (w :: 'a::len \text{ word}) = \text{bin-nth bin } (\text{size } w - 1)$   
 $\langle \text{proof} \rangle$

**lemmas**  $\text{word-msb-no} = \text{refl} [\text{THEN } \text{word-msb-no}', \text{unfolded word-size}]$

**lemma**  $\text{word-msb-nth}': \text{msb } (w :: 'a::len \text{ word}) = \text{bin-nth } (\text{uint } w) (\text{size } w - 1)$   
 $\langle \text{proof} \rangle$

**lemmas** *word-msb-nth* = *word-msb-nth'* [*unfolded word-size*]

**lemma** *word-msb-alt*: *msb* (*w*::'*a*::*len* *word*) = *hd* (*to-bl* *w*)  
 ⟨*proof*⟩

**lemma** *word-set-nth*:  
*set-bit* *w* *n* (*test-bit* *w* *n*) = (*w*::'*a*::*len0* *word*)  
 ⟨*proof*⟩

**lemma** *bin-nth-uint'*:  
*bin-nth* (*uint* *w*) *n* = (*rev* (*bin-to-bl* (*size* *w*) (*uint* *w*)) ! *n* & *n* < *size* *w*)  
 ⟨*proof*⟩

**lemmas** *bin-nth-uint* = *bin-nth-uint'* [*unfolded word-size*]

**lemma** *test-bit-bl*: *w* !! *n* = (*rev* (*to-bl* *w*) ! *n* & *n* < *size* *w*)  
 ⟨*proof*⟩

**lemma** *to-bl-nth*: *n* < *size* *w* ==> *to-bl* *w* ! *n* = *w* !! (*size* *w* - *Suc* *n*)  
 ⟨*proof*⟩

**lemma** *test-bit-set*:  
**fixes** *w* :: '*a*::*len0* *word*  
**shows** (*set-bit* *w* *n* *x*) !! *n* = (*n* < *size* *w* & *x*)  
 ⟨*proof*⟩

**lemma** *test-bit-set-gen*:  
**fixes** *w* :: '*a*::*len0* *word*  
**shows** *test-bit* (*set-bit* *w* *n* *x*) *m* =  
 (*if* *m* = *n* *then* *n* < *size* *w* & *x* *else* *test-bit* *w* *m*)  
 ⟨*proof*⟩

**lemma** *of-bl-rep-False*: *of-bl* (*replicate* *n* *False* @ *bs*) = *of-bl* *bs*  
 ⟨*proof*⟩

**lemma** *msb-nth'*:  
**fixes** *w* :: '*a*::*len* *word*  
**shows** *msb* *w* = *w* !! (*size* *w* - 1)  
 ⟨*proof*⟩

**lemmas** *msb-nth* = *msb-nth'* [*unfolded word-size*]

**lemmas** *msb0* = *len-gt-0* [*THEN* *diff-Suc-less*, *THEN*  
*word-ops-nth-size* [*unfolded word-size*], *standard*]

**lemmas** *msb1* = *msb0* [**where** *i* = 0]

**lemmas** *word-ops-msb* = *msb1* [*unfolded msb-nth* [*symmetric*, *unfolded One-nat-def*]]

**lemmas** *lsb0* = *len-gt-0* [*THEN* *word-ops-nth-size* [*unfolded word-size*], *standard*]

**lemmas** *word-ops-lsb* = *lsb0* [*unfolded word-lsb-alt*]

**lemma** *td-ext-nth'*:

*n* = *size* (*w*::'*a*::*len0* *word*) ==> *ofn* = *set-bits* ==> [*w*, *ofn* *g*] = *l* ==>  
*td-ext test-bit ofn* {*f*. *ALL i. f i --> i < n*} (%*h i. h i & i < n*)  
 <*proof*>

**lemmas** *td-ext-nth* = *td-ext-nth'* [*OF refl refl refl, unfolded word-size*]

**interpretation** *test-bit*:

*td-ext* [*op* !! :: '*a*::*len0* *word* => *nat* => *bool*  
*set-bits*  
 {*f*.  $\forall i. f i \longrightarrow i < \text{len-of TYPE('a::len0)}$ }  
 ( $\lambda h i. h i \wedge i < \text{len-of TYPE('a::len0)}$ )]  
 <*proof*>

**declare** *test-bit.Rep'* [*simp del*]

**declare** *test-bit.Rep'* [*rule del*]

**lemmas** *td-nth* = *test-bit.td-thm*

**lemma** *word-set-set-same*:

**fixes** *w* :: '*a*::*len0* *word*  
**shows** *set-bit* (*set-bit w n x*) *n y* = *set-bit w n y*  
 <*proof*>

**lemma** *word-set-set-diff*:

**fixes** *w* :: '*a*::*len0* *word*  
**assumes** *m* ~ = *n*  
**shows** *set-bit* (*set-bit w m x*) *n y* = *set-bit* (*set-bit w n y*) *m x*  
 <*proof*>

**lemma** *test-bit-no'*:

**fixes** *w* :: '*a*::*len0* *word*  
**shows** *w* = *number-of bin* ==> *test-bit w n* = (*n* < *size w* & *bin-nth bin n*)  
 <*proof*>

**lemmas** *test-bit-no* =

*refl* [*THEN test-bit-no', unfolded word-size, THEN eq-reflection, standard*]

**lemma** *nth-0*: ~ (*0*::'*a*::*len0* *word*) !! *n*

<*proof*>

**lemma** *nth-sint*:

**fixes** *w* :: '*a*::*len* *word*  
**defines** *l*  $\equiv \text{len-of TYPE ('a)}$   
**shows** *bin-nth* (*sint w*) *n* = (*if n* < *l* - 1 *then w* !! *n* *else w* !! (*l* - 1))  
 <*proof*>



**lemma** *word-lsb-no*:

*lsb (number-of bin :: 'a :: len word) = (bin-last bin = bit.B1)*  
*<proof>*

**lemma** *word-set-no*:

*set-bit (number-of bin::'a::len0 word) n b =*  
*number-of (bin-sc n (if b then bit.B1 else bit.B0) bin)*  
*<proof>*

**lemmas** *setBit-no = setBit-def [THEN trans [OF meta-eq-to-obj-eq word-set-no],*  
*simplified if-simps, THEN eq-reflection, standard]*

**lemmas** *clearBit-no = clearBit-def [THEN trans [OF meta-eq-to-obj-eq word-set-no],*  
*simplified if-simps, THEN eq-reflection, standard]*

**lemma** *to-bl-n1*:

*to-bl (-1::'a::len0 word) = replicate (len-of TYPE ('a)) True*  
*<proof>*

**lemma** *word-msb-n1*: *msb (-1::'a::len word)*

*<proof>*

**declare** *word-set-set-same [simp] word-set-nth [simp]*

*test-bit-no [simp] word-set-no [simp] nth-0 [simp]*

*setBit-no [simp] clearBit-no [simp]*

*word-lsb-no [simp] word-msb-no [simp] word-msb-n1 [simp] word-lsb-1-0 [simp]*

**lemma** *word-set-nth-iff*:

*(set-bit w n b = w) = (w !! n = b | n >= size (w::'a::len0 word))*  
*<proof>*

**lemma** *test-bit-2p'*:

*w = word-of-int (2 ^ n) ==>*

*w !! m = (m = n & m < size (w :: 'a :: len word))*

*<proof>*

**lemmas** *test-bit-2p = refl [THEN test-bit-2p', unfolded word-size]*

**lemmas** *nth-w2p = test-bit-2p [unfolded of-int-number-of-eq*  
*word-of-int [symmetric] Int.of-int-power]*

**lemma** *uint-2p*:

*(0::'a::len word) < 2 ^ n ==> uint (2 ^ n::'a::len word) = 2 ^ n*  
*<proof>*

**lemma** *word-of-int-2p*: *(word-of-int (2 ^ n) :: 'a :: len word) = 2 ^ n*  
*<proof>*

**lemma** *bang-is-le*: *x !! m ==> 2 ^ m <= (x :: 'a :: len word)*  
*<proof>*

```

lemma word-clr-le:
  fixes w :: 'a::len0 word
  shows w >= set-bit w n False
  ⟨proof⟩

```

```

lemma word-set-ge:
  fixes w :: 'a::len word
  shows w <= set-bit w n True
  ⟨proof⟩

```

```

end

```

## 13 WordShift: Shifting, Rotating, and Splitting Words

```

theory WordShift
imports WordBitwise
begin

```

### 13.1 Bit shifting

```

lemma shiffl1-number [simp] :
  shiffl1 (number-of w) = number-of (w BIT bit.B0)
  ⟨proof⟩

```

```

lemma shiffl1-0 [simp] : shiffl1 0 = 0
  ⟨proof⟩

```

```

lemmas shiffl1-def-u = shiffl1-def [folded word-number-of-def]

```

```

lemma shiffl1-def-s: shiffl1 w = number-of (sint w BIT bit.B0)
  ⟨proof⟩

```

```

lemma shiftr1-0 [simp] : shiftr1 0 = 0
  ⟨proof⟩

```

```

lemma sshiftr1-0 [simp] : sshiftr1 0 = 0
  ⟨proof⟩

```

```

lemma sshiftr1-n1 [simp] : sshiftr1 -1 = -1
  ⟨proof⟩

```

```

lemma shiffl-0 [simp] : (0::'a::len0 word) << n = 0
  ⟨proof⟩

```

**lemma** *shiftr-0* [*simp*] :  $(0::'a::len0\ word) >> n = 0$   
 ⟨*proof*⟩

**lemma** *sshiftr-0* [*simp*] :  $0 >>> n = 0$   
 ⟨*proof*⟩

**lemma** *sshiftr-n1* [*simp*] :  $-1 >>> n = -1$   
 ⟨*proof*⟩

**lemma** *nth-shiftrl1*:  $shiftrl\ w\ !!\ n = (n < size\ w \ \&\ n > 0 \ \&\ w\ !!\ (n - 1))$   
 ⟨*proof*⟩

**lemma** *nth-shiftrl'* [*rule-format*]:  
 ALL  $n. ((w::'a::len0\ word) << m) !!\ n = (n < size\ w \ \&\ n \geq m \ \&\ w\ !!\ (n - m))$   
 ⟨*proof*⟩

**lemmas** *nth-shiftrl* = *nth-shiftrl'* [*unfolded word-size*]

**lemma** *nth-shiftr1*:  $shiftr1\ w\ !!\ n = w\ !!\ Suc\ n$   
 ⟨*proof*⟩

**lemma** *nth-shiftr*:  
 $\bigwedge n. ((w::'a::len0\ word) >> m) !!\ n = w\ !!\ (n + m)$   
 ⟨*proof*⟩

**lemma** *uint-shiftr1*:  $uint\ (shiftr1\ w) = bin-rest\ (uint\ w)$   
 ⟨*proof*⟩

**lemma** *nth-sshiftr1*:  
 $sshiftr1\ w\ !!\ n = (if\ n = size\ w - 1\ then\ w\ !!\ n\ else\ w\ !!\ Suc\ n)$   
 ⟨*proof*⟩

**lemma** *nth-sshiftr* [*rule-format*] :  
 ALL  $n. sshiftr\ w\ m\ !!\ n = (n < size\ w \ \&\ (if\ n + m \geq size\ w\ then\ w\ !!\ (size\ w - 1)\ else\ w\ !!\ (n + m)))$   
 ⟨*proof*⟩

**lemma** *shiftr1-div-2*:  $uint\ (shiftr1\ w) = uint\ w\ div\ 2$   
 ⟨*proof*⟩

**lemma** *sshiftr1-div-2*:  $sint\ (sshiftr1\ w) = sint\ w\ div\ 2$   
 ⟨*proof*⟩

**lemma** *shiftr-div-2n*:  $uint\ (shiftr\ w\ n) = uint\ w\ div\ 2^{\wedge} n$   
 ⟨*proof*⟩

**lemma** *sshiftr-div-2n*:  $\text{sint} (\text{sshiftr } w \ n) = \text{sint } w \ \text{div } 2 \wedge n$   
 ⟨proof⟩

### 13.1.1 shift functions in terms of lists of bools

**lemmas** *bshiftr1-no-bin* [simp] =  
*bshiftr1-def* [where  $w = \text{number-of } w$ , *unfolded to-bl-no-bin*, *standard*]

**lemma** *bshiftr1-bl*:  $\text{to-bl} (\text{bshiftr1 } b \ w) = b \ \# \ \text{butlast} (\text{to-bl } w)$   
 ⟨proof⟩

**lemma** *shiftrl1-of-bl*:  $\text{shiftrl1} (\text{of-bl } bl) = \text{of-bl} (bl \ @ \ [\text{False}])$   
 ⟨proof⟩

**lemma** *shiftrl1-bl*:  $\text{shiftrl1} (w :: 'a :: \text{len0 word}) = \text{of-bl} (\text{to-bl } w \ @ \ [\text{False}])$   
 ⟨proof⟩

**lemma** *bl-shiftrl1*:  
 $\text{to-bl} (\text{shiftrl1} (w :: 'a :: \text{len word})) = \text{tl} (\text{to-bl } w) \ @ \ [\text{False}]$   
 ⟨proof⟩

**lemma** *shiftr1-bl*:  $\text{shiftr1 } w = \text{of-bl} (\text{butlast} (\text{to-bl } w))$   
 ⟨proof⟩

**lemma** *bl-shiftr1*:  
 $\text{to-bl} (\text{shiftr1} (w :: 'a :: \text{len word})) = \text{False} \ \# \ \text{butlast} (\text{to-bl } w)$   
 ⟨proof⟩

**lemma** *shiftrl1-rev*:  
 $\text{shiftrl1} (w :: 'a :: \text{len word}) = \text{word-reverse} (\text{shiftr1} (\text{word-reverse } w))$   
 ⟨proof⟩

**lemma** *shiftrl-rev*:  
 $\text{shiftrl} (w :: 'a :: \text{len word}) \ n = \text{word-reverse} (\text{shiftr} (\text{word-reverse } w) \ n)$   
 ⟨proof⟩

**lemmas** *rev-shiftrl* =  
*shiftrl-rev* [where  $w = \text{word-reverse } w$ , *simplified*, *standard*]

**lemmas** *shiftr-rev* = *rev-shiftrl* [THEN *word-rev-gal'*, *standard*]

**lemmas** *rev-shiftr* = *shiftrl-rev* [THEN *word-rev-gal'*, *standard*]

**lemma** *bl-sshiftr1*:  
 $\text{to-bl} (\text{sshiftr1} (w :: 'a :: \text{len word})) = \text{hd} (\text{to-bl } w) \ \# \ \text{butlast} (\text{to-bl } w)$   
 ⟨proof⟩

**lemma** *drop-shiftr*:

$\text{drop } n \text{ (to-bl ((w :: 'a :: len word) >> n))} = \text{take (size w - n) (to-bl w)}$   
 ⟨proof⟩

**lemma** *drop-sshiftr*:

$\text{drop } n \text{ (to-bl ((w :: 'a :: len word) >>> n))} = \text{take (size w - n) (to-bl w)}$   
 ⟨proof⟩

**lemma** *take-shiftr* [rule-format] :

$n \leq \text{size (w :: 'a :: len word)} \longrightarrow \text{take } n \text{ (to-bl (w >> n))} =$   
 $\text{replicate } n \text{ False}$   
 ⟨proof⟩

**lemma** *take-sshiftr'* [rule-format] :

$n \leq \text{size (w :: 'a :: len word)} \longrightarrow \text{hd (to-bl (w >>> n))} = \text{hd (to-bl w)} \ \&$   
 $\text{take } n \text{ (to-bl (w >>> n))} = \text{replicate } n \text{ (hd (to-bl w))}$   
 ⟨proof⟩

**lemmas** *hd-sshiftr* = *take-sshiftr'* [THEN conjunct1, standard]

**lemmas** *take-sshiftr* = *take-sshiftr'* [THEN conjunct2, standard]

**lemma** *atd-lem*:  $\text{take } n \text{ xs} = t \implies \text{drop } n \text{ xs} = d \implies \text{xs} = t @ d$   
 ⟨proof⟩

**lemmas** *bl-shiftr* = *atd-lem* [OF *take-shiftr drop-shiftr*]

**lemmas** *bl-sshiftr* = *atd-lem* [OF *take-sshiftr drop-sshiftr*]

**lemma** *shiftr-of-bl*:  $\text{of-bl bl} << n = \text{of-bl (bl @ replicate } n \text{ False)}$   
 ⟨proof⟩

**lemma** *shiftr-bl*:

$(w :: 'a :: \text{len0 word}) << (n :: \text{nat}) = \text{of-bl (to-bl w @ replicate } n \text{ False)}$   
 ⟨proof⟩

**lemmas** *shiftr-number* [simp] = *shiftr-def* [where *w=number-of w*, standard]

**lemma** *bl-shiftr*:

$\text{to-bl (w << n)} = \text{drop } n \text{ (to-bl w) @ replicate (min (size w) n) False}$   
 ⟨proof⟩

**lemma** *shiftr-zero-size*:

**fixes**  $x :: 'a :: \text{len0 word}$

**shows**  $\text{size } x \leq n \implies x << n = 0$

⟨proof⟩

**lemma** *shiftr1-2t*:  $\text{shiftr1 (w :: 'a :: len word)} = 2 * w$

⟨proof⟩

**lemma** *shiffl1-p*: *shiffl1* (*w* :: 'a :: len word) = *w* + *w*  
 ⟨proof⟩

**lemma** *shiffl-t2n*: *shiffl* (*w* :: 'a :: len word) *n* =  $2^{\wedge} n * w$   
 ⟨proof⟩

**lemma** *shiftr1-bintr* [simp]:  
 (*shiftr1* (number-of *w*) :: 'a :: len0 word) =  
 number-of (bin-rest (bintrunc (len-of TYPE ('a)) *w*))  
 ⟨proof⟩

**lemma** *sshiftr1-sbintr* [simp] :  
 (*sshiftr1* (number-of *w*) :: 'a :: len word) =  
 number-of (bin-rest (sbintrunc (len-of TYPE ('a) - 1) *w*))  
 ⟨proof⟩

**lemma** *shiftr-no'*:  
*w* = number-of bin ==>  
 (*w*::'a::len0 word) >> *n* = number-of ((bin-rest ^ *n*) (bintrunc (size *w*) bin))  
 ⟨proof⟩

**lemma** *sshiftr-no'*:  
*w* = number-of bin ==> *w* >>> *n* = number-of ((bin-rest ^ *n*)  
 (sbintrunc (size *w* - 1) bin))  
 ⟨proof⟩

**lemmas** *sshiftr-no* [simp] =  
*sshiftr-no'* [where *w* = number-of *w*, OF refl, unfolded word-size, standard]

**lemmas** *shiftr-no* [simp] =  
*shiftr-no'* [where *w* = number-of *w*, OF refl, unfolded word-size, standard]

**lemma** *shiftr1-bl-of'*:  
*us* = *shiftr1* (of-bl *bl*) ==> length *bl* <= size *us* ==>  
*us* = of-bl (butlast *bl*)  
 ⟨proof⟩

**lemmas** *shiftr1-bl-of* = refl [THEN *shiftr1-bl-of'*, unfolded word-size]

**lemma** *shiftr-bl-of'* [rule-format]:  
*us* = of-bl *bl* >> *n* ==> length *bl* <= size *us* -->  
*us* = of-bl (take (length *bl* - *n*) *bl*)  
 ⟨proof⟩

**lemmas** *shiftr-bl-of* = refl [THEN *shiftr-bl-of'*, unfolded word-size]

**lemmas** *shiftr-bl* = word-bl.Rep' [THEN eq-imp-le, THEN *shiftr-bl-of*,  
 simplified word-size, simplified, THEN eq-reflection, standard]

**lemma** *msb-shift'*:  $msb (w :: 'a :: len \text{ word}) <-> (w >> (size\ w - 1)) \sim = 0$   
 ⟨proof⟩

**lemmas** *msb-shift* = *msb-shift'* [unfolded word-size]

**lemma** *align-lem-or* [rule-format] :  
 ALL  $x\ m$ .  $length\ x = n + m \dashrightarrow length\ y = n + m \dashrightarrow$   
 $drop\ m\ x = replicate\ n\ False \dashrightarrow take\ m\ y = replicate\ m\ False \dashrightarrow$   
 $map2\ op\ | \ x\ y = take\ m\ x\ @\ drop\ m\ y$   
 ⟨proof⟩

**lemma** *align-lem-and* [rule-format] :  
 ALL  $x\ m$ .  $length\ x = n + m \dashrightarrow length\ y = n + m \dashrightarrow$   
 $drop\ m\ x = replicate\ n\ False \dashrightarrow take\ m\ y = replicate\ m\ False \dashrightarrow$   
 $map2\ op\ \&\ x\ y = replicate\ (n + m)\ False$   
 ⟨proof⟩

**lemma** *aligned-bl-add-size'*:  
 $size\ x - n = m \implies n \leq size\ x \implies drop\ m\ (to-bl\ x) = replicate\ n\ False$   
 $\implies$   
 $take\ m\ (to-bl\ y) = replicate\ m\ False \implies$   
 $to-bl\ (x + y) = take\ m\ (to-bl\ x)\ @\ drop\ m\ (to-bl\ y)$   
 ⟨proof⟩

**lemmas** *aligned-bl-add-size* = *reft* [THEN *aligned-bl-add-size'*]

### 13.1.2 Mask

**lemma** *nth-mask'*:  $m = mask\ n \implies test-bit\ m\ i = (i < n \ \&\ i < size\ m)$   
 ⟨proof⟩

**lemmas** *nth-mask* [simp] = *reft* [THEN *nth-mask'*]

**lemma** *mask-bl*:  $mask\ n = of-bl\ (replicate\ n\ True)$   
 ⟨proof⟩

**lemma** *mask-bin*:  $mask\ n = number-of\ (bintrunc\ n\ Int.Min)$   
 ⟨proof⟩

**lemma** *and-mask-bintr*:  $w\ AND\ mask\ n = number-of\ (bintrunc\ n\ (uint\ w))$   
 ⟨proof⟩

**lemma** *and-mask-no*:  $number-of\ i\ AND\ mask\ n = number-of\ (bintrunc\ n\ i)$   
 ⟨proof⟩

**lemmas** *and-mask-wi* = *and-mask-no* [unfolded word-number-of-def]

**lemma** *bl-and-mask*:  
 $to-bl\ (w\ AND\ mask\ n :: 'a :: len \text{ word}) =$

*replicate (len-of TYPE('a) - n) False @*  
*drop (len-of TYPE('a) - n) (to-bl w)*  
 ⟨proof⟩

**lemmas** *and-mask-mod-2p =*  
*and-mask-bintr [unfolded word-number-of-alt no-bintr-alt]*

**lemma** *and-mask-lt-2p: uint (w AND mask n) < 2 ^ n*  
 ⟨proof⟩

**lemmas** *eq-mod-iff = trans [symmetric, OF int-mod-lem eq-sym-conv]*

**lemma** *mask-eq-iff: (w AND mask n) = w <-> uint w < 2 ^ n*  
 ⟨proof⟩

**lemma** *and-mask-dvd: 2 ^ n dvd uint w = (w AND mask n = 0)*  
 ⟨proof⟩

**lemma** *and-mask-dvd-nat: 2 ^ n dvd unat w = (w AND mask n = 0)*  
 ⟨proof⟩

**lemma** *word-2p-lem:*  
*n < size w ==> w < 2 ^ n = (uint (w :: 'a :: len word) < 2 ^ n)*  
 ⟨proof⟩

**lemma** *less-mask-eq: x < 2 ^ n ==> x AND mask n = (x :: 'a :: len word)*  
 ⟨proof⟩

**lemmas** *mask-eq-iff-w2p =*  
*trans [OF mask-eq-iff word-2p-lem [symmetric], standard]*

**lemmas** *and-mask-less' =*  
*iffD2 [OF word-2p-lem and-mask-lt-2p, simplified word-size, standard]*

**lemma** *and-mask-less-size: n < size x ==> x AND mask n < 2 ^ n*  
 ⟨proof⟩

**lemma** *word-mod-2p-is-mask':*  
*c = 2 ^ n ==> c > 0 ==> x mod c = (x :: 'a :: len word) AND mask n*  
 ⟨proof⟩

**lemmas** *word-mod-2p-is-mask = refl [THEN word-mod-2p-is-mask']*

**lemma** *mask-egs:*  
*(a AND mask n) + b AND mask n = a + b AND mask n*  
*a + (b AND mask n) AND mask n = a + b AND mask n*  
*(a AND mask n) - b AND mask n = a - b AND mask n*  
*a - (b AND mask n) AND mask n = a - b AND mask n*  
*a \* (b AND mask n) AND mask n = a \* b AND mask n*



$(b \text{ AND mask } n) * a \text{ AND mask } n = b * a \text{ AND mask } n$   
 $(a \text{ AND mask } n) + (b \text{ AND mask } n) \text{ AND mask } n = a + b \text{ AND mask } n$   
 $(a \text{ AND mask } n) - (b \text{ AND mask } n) \text{ AND mask } n = a - b \text{ AND mask } n$   
 $(a \text{ AND mask } n) * (b \text{ AND mask } n) \text{ AND mask } n = a * b \text{ AND mask } n$   
 $-(a \text{ AND mask } n) \text{ AND mask } n = -a \text{ AND mask } n$   
 $\text{word-succ } (a \text{ AND mask } n) \text{ AND mask } n = \text{word-succ } a \text{ AND mask } n$   
 $\text{word-pred } (a \text{ AND mask } n) \text{ AND mask } n = \text{word-pred } a \text{ AND mask } n$   
 $\langle \text{proof} \rangle$

**lemma** *mask-power-eq*:

$(x \text{ AND mask } n) \wedge k \text{ AND mask } n = x \wedge k \text{ AND mask } n$   
 $\langle \text{proof} \rangle$

### 13.1.3 Revcast

**lemmas** *revcast-def'* = *revcast-def* [*simplified*]

**lemmas** *revcast-def''* = *revcast-def'* [*simplified word-size*]

**lemmas** *revcast-no-def* [*simp*] =

*revcast-def'* [**where** *w*=*number-of w*, *unfolded word-size*, *standard*]

**lemma** *to-bl-revcast*:

$\text{to-bl } (\text{revcast } w :: 'a :: \text{len0 word}) =$   
 $\text{takefill False } (\text{len-of TYPE } ('a)) (\text{to-bl } w)$   
 $\langle \text{proof} \rangle$

**lemma** *revcast-rev-ucast'*:

$cs = [rc, uc] ==> rc = \text{revcast } (\text{word-reverse } w) ==> uc = \text{ucast } w ==>$   
 $rc = \text{word-reverse } uc$   
 $\langle \text{proof} \rangle$

**lemmas** *revcast-rev-ucast* = *revcast-rev-ucast'* [*OF refl refl refl*]

**lemmas** *revcast-ucast* = *revcast-rev-ucast*

[**where** *w* = *word-reverse w*, *simplified word-rev-rev*, *standard*]

**lemmas** *ucast-revcast* = *revcast-rev-ucast* [*THEN word-rev-gal'*, *standard*]

**lemmas** *ucast-rev-revcast* = *revcast-ucast* [*THEN word-rev-gal'*, *standard*]

— linking revcast and cast via shift

**lemmas** *wsst-TYs* = *source-size target-size word-size*

**lemma** *revcast-down-uu'*:

$rc = \text{revcast} ==> \text{source-size } rc = \text{target-size } rc + n ==>$   
 $rc (w :: 'a :: \text{len word}) = \text{ucast } (w >> n)$   
 $\langle \text{proof} \rangle$

**lemma** *revcast-down-us'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$   
 $rc (w :: 'a :: \text{len word}) = \text{ucast } (w >>> n)$   
 $\langle \text{proof} \rangle$

**lemma** *revcast-down-su'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$   
 $rc (w :: 'a :: \text{len word}) = \text{scast } (w >> n)$   
 $\langle \text{proof} \rangle$

**lemma** *revcast-down-ss'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$   
 $rc (w :: 'a :: \text{len word}) = \text{scast } (w >>> n)$   
 $\langle \text{proof} \rangle$

**lemmas** *revcast-down-uu* = *refl* [THEN *revcast-down-uu*']

**lemmas** *revcast-down-us* = *refl* [THEN *revcast-down-us*']

**lemmas** *revcast-down-su* = *refl* [THEN *revcast-down-su*']

**lemmas** *revcast-down-ss* = *refl* [THEN *revcast-down-ss*']

**lemma** *cast-down-rev*:

$uc = \text{ucast} \implies \text{source-size } uc = \text{target-size } uc + n \implies$   
 $uc w = \text{revcast } ((w :: 'a :: \text{len word}) << n)$   
 $\langle \text{proof} \rangle$

**lemma** *revcast-up'*:

$rc = \text{revcast} \implies \text{source-size } rc + n = \text{target-size } rc \implies$   
 $rc w = (\text{ucast } w :: 'a :: \text{len word}) << n$   
 $\langle \text{proof} \rangle$

**lemmas** *revcast-up* = *refl* [THEN *revcast-up*']

**lemmas** *rc1* = *revcast-up* [THEN

*revcast-rev-ucast* [symmetric, THEN *trans*, THEN *word-rev-gal*, symmetric]]

**lemmas** *rc2* = *revcast-down-uu* [THEN

*revcast-rev-ucast* [symmetric, THEN *trans*, THEN *word-rev-gal*, symmetric]]

**lemmas** *ucast-up* =

*rc1* [simplified *rev-shiftr* [symmetric] *revcast-ucast* [symmetric]]

**lemmas** *ucast-down* =

*rc2* [simplified *rev-shiftr* *revcast-ucast* [symmetric]]

#### 13.1.4 Slices

**lemmas** *slice1-no-bin* [simp] =

*slice1-def* [where *w=number-of w*, unfolded to-bl-no-bin, standard]

**lemmas** *slice-no-bin* [simp] =

*trans* [OF *slice-def* [THEN *meta-eq-to-obj-eq*]  
*slice1-no-bin* [THEN *meta-eq-to-obj-eq*],

*unfolded word-size, standard]*

**lemma** *slice1-0* [*simp*] : *slice1 n 0 = 0*  
*<proof>*

**lemma** *slice-0* [*simp*] : *slice n 0 = 0*  
*<proof>*

**lemma** *slice-take'*: *slice n w = of-bl (take (size w - n) (to-bl w))*  
*<proof>*

**lemmas** *slice-take = slice-take'* [*unfolded word-size*]

— *shiftr* to a word of the same size is just *slice*, *slice* is just *shiftr* then *ucast*

**lemmas** *shiftr-slice = trans*  
 [*OF shiftr-bl [THEN meta-eq-to-obj-eq] slice-take [symmetric], standard]*

**lemma** *slice-shiftr*: *slice n w = ucast (w >> n)*  
*<proof>*

**lemma** *nth-slice*:  
*(slice n w :: 'a :: len0 word) !! m =*  
*(w !! (m + n) & m < len-of TYPE ('a))*  
*<proof>*

**lemma** *slice1-down-alt'*:  
*sl = slice1 n w ==> fs = size sl ==> fs + k = n ==>*  
*to-bl sl = takefill False fs (drop k (to-bl w))*  
*<proof>*

**lemma** *slice1-up-alt'*:  
*sl = slice1 n w ==> fs = size sl ==> fs = n + k ==>*  
*to-bl sl = takefill False fs (replicate k False @ (to-bl w))*  
*<proof>*

**lemmas** *sd1 = slice1-down-alt'* [*OF refl refl, unfolded word-size*]

**lemmas** *su1 = slice1-up-alt'* [*OF refl refl, unfolded word-size*]

**lemmas** *slice1-down-alt = le-add-diff-inverse [THEN sd1]*

**lemmas** *slice1-up-alt =*  
*le-add-diff-inverse [symmetric, THEN su1]*  
*le-add-diff-inverse2 [symmetric, THEN su1]*

**lemma** *ucast-slice1*: *ucast w = slice1 (size w) w*  
*<proof>*

**lemma** *ucast-slice*: *ucast w = slice 0 w*  
*<proof>*

**lemmas** *slice-id = trans [OF ucast-slice [symmetric] ucast-id]*

**lemma** *revcast-slice1'*:

$rc = \text{revcast } w \implies \text{slice1 } (\text{size } rc) \ w = rc$   
 $\langle \text{proof} \rangle$

**lemmas** *revcast-slice1* = *refl* [THEN *revcast-slice1'*]

**lemma** *slice1-tf-tf'*:

$\text{to-bl } (\text{slice1 } n \ w :: 'a :: \text{len0 word}) =$   
 $\text{rev } (\text{takefill } \text{False } (\text{len-of TYPE('a)}) (\text{rev } (\text{takefill } \text{False } n \ (\text{to-bl } w))))$   
 $\langle \text{proof} \rangle$

**lemmas** *slice1-tf-tf* = *slice1-tf-tf'*

[THEN *word-bl.Rep-inverse'*, *symmetric*, *standard*]

**lemma** *rev-slice1*:

$n + k = \text{len-of TYPE('a)} + \text{len-of TYPE('b)} \implies$   
 $\text{slice1 } n \ (\text{word-reverse } w :: 'b :: \text{len0 word}) =$   
 $\text{word-reverse } (\text{slice1 } k \ w :: 'a :: \text{len0 word})$   
 $\langle \text{proof} \rangle$

**lemma** *rev-slice'*:

$\text{res} = \text{slice } n \ (\text{word-reverse } w) \implies n + k + \text{size } \text{res} = \text{size } w \implies$   
 $\text{res} = \text{word-reverse } (\text{slice } k \ w)$   
 $\langle \text{proof} \rangle$

**lemmas** *rev-slice* = *refl* [THEN *rev-slice'*, *unfolded word-size*]

**lemmas** *sym-notr* =

*not-iff* [THEN *iffD2*, THEN *not-sym*, THEN *not-iff* [THEN *iffD1*]]

— problem posed by TPHOLs referee: criterion for overflow of addition of signed integers

**lemma** *soft-test*:

$(\text{sint } (x :: 'a :: \text{len word}) + \text{sint } y = \text{sint } (x + y)) =$   
 $((((x+y) \text{ XOR } x) \text{ AND } ((x+y) \text{ XOR } y)) >> (\text{size } x - 1) = 0)$   
 $\langle \text{proof} \rangle$

## 13.2 Split and cat

**lemmas** *word-split-bin'* = *word-split-def* [THEN *meta-eq-to-obj-eq*, *standard*]

**lemmas** *word-cat-bin'* = *word-cat-def* [THEN *meta-eq-to-obj-eq*, *standard*]

**lemma** *word-rsplit-no*:

$(\text{word-rsplit } (\text{number-of bin} :: 'b :: \text{len0 word}) :: 'a \text{ word list}) =$   
 $\text{map number-of } (\text{bin-rsplit } (\text{len-of TYPE('a)} :: \text{len}))$   
 $(\text{len-of TYPE('b)}, \text{bintrunc } (\text{len-of TYPE('b)}) \text{ bin}))$   
 $\langle \text{proof} \rangle$

**lemmas** *word-rsplit-no-cl* [*simp*] = *word-rsplit-no*  
 [unfolded *bin-rsplittl-def* *bin-rsplit-l* [*symmetric*]]

**lemma** *test-bit-cat*:

$wc = \text{word-cat } a \ b ==> \text{wc} !! n = (n < \text{size } wc \ \& \$   
 $(\text{if } n < \text{size } b \text{ then } b !! n \text{ else } a !! (n - \text{size } b)))$   
 $\langle \text{proof} \rangle$

**lemma** *word-cat-bl*:  $\text{word-cat } a \ b = \text{of-bl } (\text{to-bl } a \ @ \ \text{to-bl } b)$   
 $\langle \text{proof} \rangle$

**lemma** *of-bl-append*:

$(\text{of-bl } (xs \ @ \ ys) :: 'a :: \text{len word}) = \text{of-bl } xs * 2^{(\text{length } ys)} + \text{of-bl } ys$   
 $\langle \text{proof} \rangle$

**lemma** *of-bl-False* [*simp*]:

$\text{of-bl } (\text{False} \# xs) = \text{of-bl } xs$   
 $\langle \text{proof} \rangle$

**lemma** *of-bl-True*:

$(\text{of-bl } (\text{True} \# xs) :: 'a :: \text{len word}) = 2^{\text{length } xs} + \text{of-bl } xs$   
 $\langle \text{proof} \rangle$

**lemma** *of-bl-Cons*:

$\text{of-bl } (x \# xs) = \text{of-bool } x * 2^{\text{length } xs} + \text{of-bl } xs$   
 $\langle \text{proof} \rangle$

**lemma** *split-uint-lem*:  $\text{bin-split } n \ (\text{uint } (w :: 'a :: \text{len0 word})) = (a, b) ==>$   
 $a = \text{bintrunc } (\text{len-of TYPE('a)} - n) \ a \ \& \ b = \text{bintrunc } (\text{len-of TYPE('a)}) \ b$   
 $\langle \text{proof} \rangle$

**lemma** *word-split-bl'*:

$\text{std} = \text{size } c - \text{size } b ==> (\text{word-split } c = (a, b)) ==>$   
 $(a = \text{of-bl } (\text{take } \text{std } (\text{to-bl } c)) \ \& \ b = \text{of-bl } (\text{drop } \text{std } (\text{to-bl } c)))$   
 $\langle \text{proof} \rangle$

**lemma** *word-split-bl*:  $\text{std} = \text{size } c - \text{size } b ==>$

$(a = \text{of-bl } (\text{take } \text{std } (\text{to-bl } c)) \ \& \ b = \text{of-bl } (\text{drop } \text{std } (\text{to-bl } c))) <->$   
 $\text{word-split } c = (a, b)$   
 $\langle \text{proof} \rangle$

**lemma** *word-split-bl-eq*:

$(\text{word-split } (c :: 'a :: \text{len word}) :: ('c :: \text{len0 word} * 'd :: \text{len0 word})) =$   
 $(\text{of-bl } (\text{take } (\text{len-of TYPE('a)} - \text{len-of TYPE('d)})) (\text{to-bl } c)),$   
 $\text{of-bl } (\text{drop } (\text{len-of TYPE('a)} - \text{len-of TYPE('d)}) (\text{to-bl } c)))$   
 $\langle \text{proof} \rangle$

**lemma** *test-bit-split'*:

$\text{word-split } c = (a, b) --> (\text{ALL } n \ m. \ b !! n = (n < \text{size } b \ \& \ c !! n) \ \& \$

$a !! m = (m < \text{size } a \ \& \ c !! (m + \text{size } b))$   
 $\langle \text{proof} \rangle$

**lemmas** *test-bit-split* =  
*test-bit-split'* [THEN *mp*, *simplified all-simps*, *standard*]

**lemma** *test-bit-split-eq*:  $\text{word-split } c = (a, b) <->$   
 $((\text{ALL } n::\text{nat}. b !! n = (n < \text{size } b \ \& \ c !! n)) \ \& \\$   
 $(\text{ALL } m::\text{nat}. a !! m = (m < \text{size } a \ \& \ c !! (m + \text{size } b))))$   
 $\langle \text{proof} \rangle$

**lemma** *word-cat-id*:  $\text{word-cat } a \ b = b$   
 $\langle \text{proof} \rangle$

**lemma** *word-cat-hom*:  
 $\text{len-of TYPE}('a::\text{len0}) \leq \text{len-of TYPE}('b::\text{len0}) + \text{len-of TYPE}('c::\text{len0})$   
 $==>$   
 $(\text{word-cat } (\text{word-of-int } w :: 'b \text{ word}) (b :: 'c \text{ word}) :: 'a \text{ word}) =$   
 $\text{word-of-int } (\text{bin-cat } w (\text{size } b) (\text{uint } b))$   
 $\langle \text{proof} \rangle$

**lemma** *word-cat-split-alt*:  
 $\text{size } w \leq \text{size } u + \text{size } v ==> \text{word-split } w = (u, v) ==> \text{word-cat } u \ v = w$   
 $\langle \text{proof} \rangle$

**lemmas** *word-cat-split-size* =  
*sym* [THEN [2] *word-cat-split-alt* [*symmetric*], *standard*]

### 13.2.1 Split and slice

**lemma** *split-slices*:  
 $\text{word-split } w = (u, v) ==> u = \text{slice } (\text{size } v) \ w \ \& \ v = \text{slice } 0 \ w$   
 $\langle \text{proof} \rangle$

**lemma** *slice-cat1'*:  
 $wc = \text{word-cat } a \ b ==> \text{size } wc \geq \text{size } a + \text{size } b ==> \text{slice } (\text{size } b) \ wc = a$   
 $\langle \text{proof} \rangle$

**lemmas** *slice-cat1* = *refl* [THEN *slice-cat1'*]  
**lemmas** *slice-cat2* = *trans* [OF *slice-id* *word-cat-id*]

**lemma** *cat-slices*:  
 $a = \text{slice } n \ c ==> b = \text{slice } 0 \ c ==> n = \text{size } b ==>$   
 $\text{size } a + \text{size } b \geq \text{size } c ==> \text{word-cat } a \ b = c$   
 $\langle \text{proof} \rangle$

**lemma** *word-split-cat-alt*:  
 $w = \text{word-cat } u \ v ==> \text{size } u + \text{size } v \leq \text{size } w ==> \text{word-split } w = (u, v)$   
 $\langle \text{proof} \rangle$

**lemmas** *word-cat-bl-no-bin* [simp] =  
*word-cat-bl* [where *a=number-of a*  
 and *b=number-of b*,  
 unfolded to-bl-no-bin, standard]

**lemmas** *word-split-bl-no-bin* [simp] =  
*word-split-bl-eq* [where *c=number-of c*, unfolded to-bl-no-bin, standard]

— this odd result arises from the fact that the statement of the result implies that the decoded words are of the same type, and therefore of the same length, as the original word

**lemma** *word-rsplit-same*: *word-rsplit w* = [*w*]  
 ⟨proof⟩

**lemma** *word-rsplit-empty-iff-size*:  
 (*word-rsplit w* = []) = (*size w* = 0)  
 ⟨proof⟩

**lemma** *test-bit-rsplit*:  
*sw* = *word-rsplit w* ==> *m* < *size (hd sw :: 'a :: len word)* ==>  
*k* < *length sw* ==> (*rev sw* ! *k*) !! *m* = (*w* !! (*k* \* *size (hd sw)* + *m*))  
 ⟨proof⟩

**lemma** *word-rcat-bl*: *word-rcat wl* == *of-bl (concat (map to-bl wl))*  
 ⟨proof⟩

**lemma** *size-rcat-lem'*:  
*size (concat (map to-bl wl))* = *length wl* \* *size (hd wl)*  
 ⟨proof⟩

**lemmas** *size-rcat-lem* = *size-rcat-lem'* [unfolded word-size]

**lemmas** *td-gal-lt-len* = *len-gt-0* [THEN *td-gal-lt*, standard]

**lemma** *nth-rcat-lem'* [rule-format] :  
*sw* = *size (hd wl :: 'a :: len word)* ==> (ALL *n*. *n* < *size wl* \* *sw* -->  
*rev (concat (map to-bl wl))* ! *n* =  
*rev (to-bl (rev wl ! (n div sw)))* ! (*n mod sw*))  
 ⟨proof⟩

**lemmas** *nth-rcat-lem* = *refl* [THEN *nth-rcat-lem'*, unfolded word-size]

**lemma** *test-bit-rcat*:  
*sw* = *size (hd wl :: 'a :: len word)* ==> *rc* = *word-rcat wl* ==> *rc* !! *n* =  
 (*n* < *size rc* & *n div sw* < *size wl* & (*rev wl* ! (*n div sw*) !! (*n mod sw*))  
 ⟨proof⟩

**lemma** *foldl-eq-foldr* [rule-format] :

*ALL*  $x$ .  $\text{foldl } op + x \, xs = \text{foldr } op + (x \, \# \, xs) \, (0 :: 'a :: \text{comm-monoid-add})$   
 ⟨proof⟩

**lemmas**  $\text{test-bit-cong} = \text{arg-cong} \, [\text{where } f = \text{test-bit}, \text{ THEN } \text{fun-cong}]$

**lemmas**  $\text{test-bit-rsplit-alt} =$   
 $\text{trans} \, [OF \, \text{nth-rev-alt} \, [THEN \, \text{test-bit-cong}]$   
 $\text{test-bit-rsplit} \, [OF \, \text{refl asm-rl diff-Suc-less}]]$

— lazy way of expressing that  $u$  and  $v$ , and  $su$  and  $sv$ , have same types

**lemma**  $\text{word-rsplit-len-indep}'$ :  
 $[u, v] = p ==> [su, sv] = q ==> \text{word-rsplit } u = su ==>$   
 $\text{word-rsplit } v = sv ==> \text{length } su = \text{length } sv$   
 ⟨proof⟩

**lemmas**  $\text{word-rsplit-len-indep} = \text{word-rsplit-len-indep}' \, [OF \, \text{refl refl refl refl}]$

**lemma**  $\text{length-word-rsplit-size}$ :  
 $n = \text{len-of } TYPE \, ('a :: \text{len}) ==>$   
 $(\text{length } (\text{word-rsplit } w :: 'a \text{ word list}) \leq m) = (\text{size } w \leq m * n)$   
 ⟨proof⟩

**lemmas**  $\text{length-word-rsplit-lt-size} =$   
 $\text{length-word-rsplit-size} \, [\text{unfolded } \text{Not-eq-iff linorder-not-less} \, [\text{symmetric}]]$

**lemma**  $\text{length-word-rsplit-exp-size}$ :  
 $n = \text{len-of } TYPE \, ('a :: \text{len}) ==>$   
 $\text{length } (\text{word-rsplit } w :: 'a \text{ word list}) = (\text{size } w + n - 1) \text{ div } n$   
 ⟨proof⟩

**lemma**  $\text{length-word-rsplit-even-size}$ :  
 $n = \text{len-of } TYPE \, ('a :: \text{len}) ==> \text{size } w = m * n ==>$   
 $\text{length } (\text{word-rsplit } w :: 'a \text{ word list}) = m$   
 ⟨proof⟩

**lemmas**  $\text{length-word-rsplit-exp-size}' = \text{refl} \, [THEN \, \text{length-word-rsplit-exp-size}]$

**lemmas**  $\text{tdle} = \text{iffD2} \, [OF \, \text{split-div-lemma refl}, \text{ THEN } \text{conjunct1}]$

**lemmas**  $\text{dtle} = \text{xtr4} \, [OF \, \text{tdle mult-commute}]$

**lemma**  $\text{word-rcat-rsplit}$ :  $\text{word-rcat } (\text{word-rsplit } w) = w$   
 ⟨proof⟩

**lemma**  $\text{size-word-rsplit-rcat-size}'$ :  
 $\text{word-rcat } (ws :: 'a :: \text{len word list}) = \text{frcw} ==>$   
 $\text{size } \text{frcw} = \text{length } ws * \text{len-of } TYPE \, ('a) ==>$   
 $\text{size } (\text{hd } [\text{word-rsplit } \text{frcw}, ws]) = \text{size } ws$   
 ⟨proof⟩



**lemmas** *size-word-rsplit-rcat-size* =  
*size-word-rsplit-rcat-size'* [*simplified*]

**lemma** *msreus*:  
**fixes** *n::nat*  
**shows**  $0 < n \implies (k * n + m) \text{ div } n = m \text{ div } n + k$   
**and**  $(k * n + m) \text{ mod } n = m \text{ mod } n$   
 $\langle \text{proof} \rangle$

**lemma** *word-rsplit-rcat-size'*:  
 $\text{word-rcat } (ws :: 'a :: \text{len word list}) = \text{frcw} \implies$   
 $\text{size frcw} = \text{length } ws * \text{len-of TYPE } ('a) \implies \text{word-rsplit frcw} = ws$   
 $\langle \text{proof} \rangle$

**lemmas** *word-rsplit-rcat-size* = *refl* [*THEN word-rsplit-rcat-size'*]

### 13.3 Rotation

**lemmas** *rotater-0'* [*simp*] = *rotater-def* [**where** *n* = 0, *simplified*]

**lemmas** *word-rot-defs* = *word-roti-def word-rotr-def word-rotl-def*

**lemma** *rotate-eq-mod*:  
 $m \text{ mod length } xs = n \text{ mod length } xs \implies \text{rotate } m \text{ } xs = \text{rotate } n \text{ } xs$   
 $\langle \text{proof} \rangle$

**lemmas** *rotate-eqs* [*standard*] =  
*trans* [*OF rotate0* [*THEN fun-cong*] *id-apply*]  
*rotate-rotate* [*symmetric*]  
*rotate-id*  
*rotate-conv-mod*  
*rotate-eq-mod*

#### 13.3.1 Rotation of list to right

**lemma** *rotate1-rl'*:  $\text{rotater1 } (l @ [a]) = a \# l$   
 $\langle \text{proof} \rangle$

**lemma** *rotate1-rl* [*simp*] :  $\text{rotater1 } (\text{rotate1 } l) = l$   
 $\langle \text{proof} \rangle$

**lemma** *rotate1-lr* [*simp*] :  $\text{rotate1 } (\text{rotater1 } l) = l$   
 $\langle \text{proof} \rangle$

**lemma** *rotater1-rev'*:  $\text{rotater1 } (\text{rev } xs) = \text{rev } (\text{rotate1 } xs)$   
 $\langle \text{proof} \rangle$

**lemma** *rotater-rev'*:  $\text{rotater } n \text{ } (\text{rev } xs) = \text{rev } (\text{rotate } n \text{ } xs)$   
 $\langle \text{proof} \rangle$

**lemmas** *rotater-rev* = *rotater-rev'* [where *xs* = *rev ys*, *simplified*, *standard*]

**lemma** *rotater-drop-take*:

*rotater* *n xs* =  
   *drop* (*length xs* - *n mod length xs*) *xs* @  
   *take* (*length xs* - *n mod length xs*) *xs*  
 ⟨*proof*⟩

**lemma** *rotater-Suc* [*simp*] :

*rotater* (*Suc n*) *xs* = *rotater1* (*rotater n xs*)  
 ⟨*proof*⟩

**lemma** *rotate-inv-plus* [*rule-format*] :

ALL *k*. *k* = *m* + *n*  $\rightarrow$  *rotater k* (*rotate n xs*) = *rotater m xs* &  
   *rotate k* (*rotater n xs*) = *rotate m xs* &  
   *rotater n* (*rotate k xs*) = *rotate m xs* &  
   *rotate n* (*rotater k xs*) = *rotater m xs*  
 ⟨*proof*⟩

**lemmas** *rotate-inv-rel* = *le-add-diff-inverse2* [*symmetric*, *THEN rotate-inv-plus*]

**lemmas** *rotate-inv-eq* = *order-refl* [*THEN rotate-inv-rel*, *simplified*]

**lemmas** *rotate-lr* [*simp*] = *rotate-inv-eq* [*THEN conjunct1*, *standard*]

**lemmas** *rotate-rl* [*simp*] =  
   *rotate-inv-eq* [*THEN conjunct2*, *THEN conjunct1*, *standard*]

**lemma** *rotate-gal*: (*rotater n xs* = *ys*) = (*rotate n ys* = *xs*)  
 ⟨*proof*⟩

**lemma** *rotate-gal'*: (*ys* = *rotater n xs*) = (*xs* = *rotate n ys*)  
 ⟨*proof*⟩

**lemma** *length-rotater* [*simp*]:

*length* (*rotater n xs*) = *length xs*  
 ⟨*proof*⟩

**lemmas** *rrs0* = *rotate-eqs* [*THEN restrict-to-left*,  
   *simplified rotate-gal* [*symmetric*] *rotate-gal'* [*symmetric*], *standard*]

**lemmas** *rrs1* = *rrs0* [*THEN refl* [*THEN rev-iffD1*]]

**lemmas** *rotater-eqs* = *rrs1* [*simplified length-rotater*, *standard*]

**lemmas** *rotater-0* = *rotater-eqs* (1)

**lemmas** *rotater-add* = *rotater-eqs* (2)

### 13.3.2 map, map2, commuting with rotate(r)

**lemma** *last-map*: *xs*  $\sim$  []  $\Rightarrow$  *last* (*map f xs*) = *f* (*last xs*)  
 ⟨*proof*⟩

**lemma** *butlast-map*:

$xs \sim [] \implies \text{butlast } (\text{map } f \text{ } xs) = \text{map } f \text{ } (\text{butlast } xs)$   
 ⟨proof⟩

**lemma** *rotater1-map*:  $\text{rotater1 } (\text{map } f \text{ } xs) = \text{map } f \text{ } (\text{rotater1 } xs)$

⟨proof⟩

**lemma** *rotater-map*:

$\text{rotater } n \text{ } (\text{map } f \text{ } xs) = \text{map } f \text{ } (\text{rotater } n \text{ } xs)$   
 ⟨proof⟩

**lemma** *but-last-zip* [rule-format] :

$ALL \text{ } ys. \text{length } xs = \text{length } ys \implies xs \sim [] \implies$   
 $\text{last } (\text{zip } xs \text{ } ys) = (\text{last } xs, \text{last } ys) \ \&$   
 $\text{butlast } (\text{zip } xs \text{ } ys) = \text{zip } (\text{butlast } xs) \text{ } (\text{butlast } ys)$   
 ⟨proof⟩

**lemma** *but-last-map2* [rule-format] :

$ALL \text{ } ys. \text{length } xs = \text{length } ys \implies xs \sim [] \implies$   
 $\text{last } (\text{map2 } f \text{ } xs \text{ } ys) = f \text{ } (\text{last } xs) \text{ } (\text{last } ys) \ \&$   
 $\text{butlast } (\text{map2 } f \text{ } xs \text{ } ys) = \text{map2 } f \text{ } (\text{butlast } xs) \text{ } (\text{butlast } ys)$   
 ⟨proof⟩

**lemma** *rotater1-zip*:

$\text{length } xs = \text{length } ys \implies$   
 $\text{rotater1 } (\text{zip } xs \text{ } ys) = \text{zip } (\text{rotater1 } xs) \text{ } (\text{rotater1 } ys)$   
 ⟨proof⟩

**lemma** *rotater1-map2*:

$\text{length } xs = \text{length } ys \implies$   
 $\text{rotater1 } (\text{map2 } f \text{ } xs \text{ } ys) = \text{map2 } f \text{ } (\text{rotater1 } xs) \text{ } (\text{rotater1 } ys)$   
 ⟨proof⟩

**lemmas** *lrth* =

$\text{box-equals } [OF \text{ } \text{asm-rl } \text{length-rotater } [symmetric]$   
 $\text{length-rotater } [symmetric],$   
 $THEN \text{ } \text{rotater1-map2}]$

**lemma** *rotater-map2*:

$\text{length } xs = \text{length } ys \implies$   
 $\text{rotater } n \text{ } (\text{map2 } f \text{ } xs \text{ } ys) = \text{map2 } f \text{ } (\text{rotater } n \text{ } xs) \text{ } (\text{rotater } n \text{ } ys)$   
 ⟨proof⟩

**lemma** *rotate1-map2*:

$\text{length } xs = \text{length } ys \implies$   
 $\text{rotate1 } (\text{map2 } f \text{ } xs \text{ } ys) = \text{map2 } f \text{ } (\text{rotate1 } xs) \text{ } (\text{rotate1 } ys)$   
 ⟨proof⟩

**lemmas** *lth = box-equals* [*OF asm-rl length-rotate* [*symmetric*]  
*length-rotate* [*symmetric*], *THEN rotate1-map2*]

**lemma** *rotate-map2*:  
*length xs = length ys ==>*  
*rotate n (map2 f xs ys) = map2 f (rotate n xs) (rotate n ys)*  
*<proof>*

**lemma** *to-bl-rotl*:  
*to-bl (word-rotl n w) = rotate n (to-bl w)*  
*<proof>*

**lemmas** *blrs0 = rotate-egs* [*THEN to-bl-rotl* [*THEN trans*]]

**lemmas** *word-rotl-egs =*  
*blrs0* [*simplified word-bl.Rep' word-bl.Rep-inject to-bl-rotl* [*symmetric*]]

**lemma** *to-bl-rotr*:  
*to-bl (word-rotr n w) = rotater n (to-bl w)*  
*<proof>*

**lemmas** *brrs0 = rotater-egs* [*THEN to-bl-rotr* [*THEN trans*]]

**lemmas** *word-rotr-egs =*  
*brrs0* [*simplified word-bl.Rep' word-bl.Rep-inject to-bl-rotr* [*symmetric*]]

**declare** *word-rotr-egs* (1) [*simp*]  
**declare** *word-rotl-egs* (1) [*simp*]

**lemma**  
*word-rot-rl* [*simp*]:  
*word-rotl k (word-rotr k v) = v and*  
*word-rot-lr* [*simp*]:  
*word-rotr k (word-rotl k v) = v*  
*<proof>*

**lemma**  
*word-rot-gal*:  
*(word-rotr n v = w) = (word-rotl n w = v) and*  
*word-rot-gal'*:  
*(w = word-rotr n v) = (v = word-rotl n w)*  
*<proof>*

**lemma** *word-rotr-rev*:  
*word-rotr n w = word-reverse (word-rotl n (word-reverse w))*  
*<proof>*

**lemma** *word-roti-0* [*simp*]: *word-roti 0 w = w*  
*<proof>*

**lemmas** *abl-cong* = *arg-cong* [where *f* = *of-bl*]

**lemma** *word-roti-add*:

*word-roti* (*m* + *n*) *w* = *word-roti* *m* (*word-roti* *n* *w*)  
 ⟨*proof*⟩

**lemma** *word-roti-conv-mod'*: *word-roti* *n* *w* = *word-roti* (*n mod int (size w)*) *w*  
 ⟨*proof*⟩

**lemmas** *word-roti-conv-mod* = *word-roti-conv-mod'* [unfolded *word-size*]

### 13.3.3 Word rotation commutes with bit-wise operations

**locale** *word-rotate*

**context** *word-rotate*

**begin**

**lemmas** *word-rot-defs'* = *to-bl-rotl to-bl-rotr*

**lemmas** *blwl-syms* [symmetric] = *bl-word-not bl-word-and bl-word-or bl-word-xor*

**lemmas** *lbl-lbl* = *trans* [OF *word-bl.Rep'* *word-bl.Rep'* [symmetric]]

**lemmas** *ths-map2* [OF *lbl-lbl*] = *rotate-map2 rotater-map2*

**lemmas** *ths-map* [where *xs* = *to-bl v*, *standard*] = *rotate-map rotater-map*

**lemmas** *th1s* [simplified *word-rot-defs'* [symmetric]] = *ths-map2 ths-map*

**lemma** *word-rot-logs*:

*word-rotl* *n* (*NOT v*) = *NOT word-rotl* *n* *v*  
*word-rotr* *n* (*NOT v*) = *NOT word-rotr* *n* *v*  
*word-rotl* *n* (*x AND y*) = *word-rotl* *n* *x AND word-rotl* *n* *y*  
*word-rotr* *n* (*x AND y*) = *word-rotr* *n* *x AND word-rotr* *n* *y*  
*word-rotl* *n* (*x OR y*) = *word-rotl* *n* *x OR word-rotl* *n* *y*  
*word-rotr* *n* (*x OR y*) = *word-rotr* *n* *x OR word-rotr* *n* *y*  
*word-rotl* *n* (*x XOR y*) = *word-rotl* *n* *x XOR word-rotl* *n* *y*  
*word-rotr* *n* (*x XOR y*) = *word-rotr* *n* *x XOR word-rotr* *n* *y*  
 ⟨*proof*⟩

**end**

**lemmas** *word-rot-logs* = *word-rotate.word-rot-logs*

**lemmas** *bl-word-rotl-dt* = *trans* [OF *to-bl-rotl rotate-drop-take*,  
*simplified word-bl.Rep'*, *standard*]

**lemmas** *bl-word-rotr-dt* = *trans* [OF *to-bl-rotr rotater-drop-take*,

```

simplified word-bl.Rep', standard]

lemma bl-word-roti-dt':
  n = nat ((- i) mod int (size (w :: 'a :: len word))) ==>
    to-bl (word-roti i w) = drop n (to-bl w) @ take n (to-bl w)
  <proof>

lemmas bl-word-roti-dt = bl-word-roti-dt' [unfolded word-size]

lemmas word-rotl-dt = bl-word-rotl-dt
  [THEN word-bl.Rep-inverse' [symmetric], standard]
lemmas word-rotr-dt = bl-word-rotr-dt
  [THEN word-bl.Rep-inverse' [symmetric], standard]
lemmas word-roti-dt = bl-word-roti-dt
  [THEN word-bl.Rep-inverse' [symmetric], standard]

lemma word-rotx-0 [simp] : word-rotr i 0 = 0 & word-rotl i 0 = 0
  <proof>

lemma word-roti-0' [simp] : word-roti n 0 = 0
  <proof>

lemmas word-rotr-dt-no-bin' [simp] =
  word-rotr-dt [where w=number-of w, unfolded to-bl-no-bin, standard]

lemmas word-rotl-dt-no-bin' [simp] =
  word-rotl-dt [where w=number-of w, unfolded to-bl-no-bin, standard]

declare word-roti-def [simp]

end

```

## 14 Boolean-Algebra: Boolean Algebras

```

theory Boolean-Algebra
imports ATP-Linkup
begin

locale boolean =
  fixes conj :: 'a ⇒ 'a ⇒ 'a (infixr □ 70)
  fixes disj :: 'a ⇒ 'a ⇒ 'a (infixr ⊔ 65)
  fixes compl :: 'a ⇒ 'a (∼ - [81] 80)
  fixes zero :: 'a (0)
  fixes one  :: 'a (1)
  assumes conj-assoc: (x □ y) □ z = x □ (y □ z)
  assumes disj-assoc: (x ⊔ y) ⊔ z = x ⊔ (y ⊔ z)
  assumes conj-commute: x □ y = y □ x

```

**assumes** *disj-commute*:  $x \sqcup y = y \sqcup x$   
**assumes** *conj-disj-distrib*:  $x \sqcap (y \sqcup z) = (x \sqcap y) \sqcup (x \sqcap z)$   
**assumes** *disj-conj-distrib*:  $x \sqcup (y \sqcap z) = (x \sqcup y) \sqcap (x \sqcup z)$   
**assumes** *conj-one-right* [*simp*]:  $x \sqcap \mathbf{1} = x$   
**assumes** *disj-zero-right* [*simp*]:  $x \sqcup \mathbf{0} = x$   
**assumes** *conj-cancel-right* [*simp*]:  $x \sqcap \sim x = \mathbf{0}$   
**assumes** *disj-cancel-right* [*simp*]:  $x \sqcup \sim x = \mathbf{1}$   
**begin**

**lemmas** *disj-ac* =  
*disj-assoc disj-commute*  
*mk-left-commute* [**where**  $'a = 'a$ , *of disj*, *OF disj-assoc disj-commute*]

**lemmas** *conj-ac* =  
*conj-assoc conj-commute*  
*mk-left-commute* [**where**  $'a = 'a$ , *of conj*, *OF conj-assoc conj-commute*]

**lemma** *dual*: *boolean disj conj compl one zero*  
*<proof>*

## 14.1 Complement

**lemma** *complement-unique*:  
**assumes** *1*:  $a \sqcap x = \mathbf{0}$   
**assumes** *2*:  $a \sqcup x = \mathbf{1}$   
**assumes** *3*:  $a \sqcap y = \mathbf{0}$   
**assumes** *4*:  $a \sqcup y = \mathbf{1}$   
**shows**  $x = y$   
*<proof>*

**lemma** *compl-unique*:  $\llbracket x \sqcap y = \mathbf{0}; x \sqcup y = \mathbf{1} \rrbracket \implies \sim x = y$   
*<proof>*

**lemma** *double-compl* [*simp*]:  $\sim(\sim x) = x$   
*<proof>*

**lemma** *compl-eq-compl-iff* [*simp*]:  $(\sim x = \sim y) = (x = y)$   
*<proof>*

## 14.2 Conjunction

**lemma** *conj-absorb* [*simp*]:  $x \sqcap x = x$   
*<proof>*

**lemma** *conj-zero-right* [*simp*]:  $x \sqcap \mathbf{0} = \mathbf{0}$   
*<proof>*

**lemma** *compl-one* [*simp*]:  $\sim \mathbf{1} = \mathbf{0}$   
*<proof>*

**lemma** *conj-zero-left* [*simp*]:  $\mathbf{0} \sqcap x = \mathbf{0}$   
 $\langle proof \rangle$

**lemma** *conj-one-left* [*simp*]:  $\mathbf{1} \sqcap x = x$   
 $\langle proof \rangle$

**lemma** *conj-cancel-left* [*simp*]:  $\sim x \sqcap x = \mathbf{0}$   
 $\langle proof \rangle$

**lemma** *conj-left-absorb* [*simp*]:  $x \sqcap (x \sqcap y) = x \sqcap y$   
 $\langle proof \rangle$

**lemma** *conj-disj-distrib2*:  
 $(y \sqcup z) \sqcap x = (y \sqcap x) \sqcup (z \sqcap x)$   
 $\langle proof \rangle$

**lemmas** *conj-disj-distrib* =  
*conj-disj-distrib conj-disj-distrib2*

### 14.3 Disjunction

**lemma** *disj-absorb* [*simp*]:  $x \sqcup x = x$   
 $\langle proof \rangle$

**lemma** *disj-one-right* [*simp*]:  $x \sqcup \mathbf{1} = \mathbf{1}$   
 $\langle proof \rangle$

**lemma** *compl-zero* [*simp*]:  $\sim \mathbf{0} = \mathbf{1}$   
 $\langle proof \rangle$

**lemma** *disj-zero-left* [*simp*]:  $\mathbf{0} \sqcup x = x$   
 $\langle proof \rangle$

**lemma** *disj-one-left* [*simp*]:  $\mathbf{1} \sqcup x = \mathbf{1}$   
 $\langle proof \rangle$

**lemma** *disj-cancel-left* [*simp*]:  $\sim x \sqcup x = \mathbf{1}$   
 $\langle proof \rangle$

**lemma** *disj-left-absorb* [*simp*]:  $x \sqcup (x \sqcup y) = x \sqcup y$   
 $\langle proof \rangle$

**lemma** *disj-conj-distrib2*:  
 $(y \sqcap z) \sqcup x = (y \sqcup x) \sqcap (z \sqcup x)$   
 $\langle proof \rangle$

**lemmas** *disj-conj-distrib* =  
*disj-conj-distrib disj-conj-distrib2*



### 14.4 De Morgan’s Laws

**lemma** *de-Morgan-conj* [*simp*]:  $\sim (x \sqcap y) = \sim x \sqcup \sim y$   
 $\langle \text{proof} \rangle$

**lemma** *de-Morgan-disj* [*simp*]:  $\sim (x \sqcup y) = \sim x \sqcap \sim y$   
 $\langle \text{proof} \rangle$

**end**

### 14.5 Symmetric Difference

**locale** *boolean-xor* = *boolean* +  
**fixes** *xor* :: 'a => 'a => 'a (**infixr**  $\oplus$  65)  
**assumes** *xor-def*:  $x \oplus y = (x \sqcap \sim y) \sqcup (\sim x \sqcap y)$   
**begin**

**lemma** *xor-def2*:  
 $x \oplus y = (x \sqcup y) \sqcap (\sim x \sqcup \sim y)$   
 $\langle \text{proof} \rangle$

**lemma** *xor-commute*:  $x \oplus y = y \oplus x$   
 $\langle \text{proof} \rangle$

**lemma** *xor-assoc*:  $(x \oplus y) \oplus z = x \oplus (y \oplus z)$   
 $\langle \text{proof} \rangle$

**lemmas** *xor-ac* =  
*xor-assoc xor-commute*  
*mk-left-commute* [**where** 'a = 'a, of *xor*, OF *xor-assoc xor-commute*]

**lemma** *xor-zero-right* [*simp*]:  $x \oplus \mathbf{0} = x$   
 $\langle \text{proof} \rangle$

**lemma** *xor-zero-left* [*simp*]:  $\mathbf{0} \oplus x = x$   
 $\langle \text{proof} \rangle$

**lemma** *xor-one-right* [*simp*]:  $x \oplus \mathbf{1} = \sim x$   
 $\langle \text{proof} \rangle$

**lemma** *xor-one-left* [*simp*]:  $\mathbf{1} \oplus x = \sim x$   
 $\langle \text{proof} \rangle$

**lemma** *xor-self* [*simp*]:  $x \oplus x = \mathbf{0}$   
 $\langle \text{proof} \rangle$

**lemma** *xor-left-self* [*simp*]:  $x \oplus (x \oplus y) = y$   
 $\langle \text{proof} \rangle$

**lemma** *xor-compl-left*:  $\sim x \oplus y = \sim (x \oplus y)$

$\langle proof \rangle$

**lemma** *xor-compl-right*:  $x \oplus \sim y = \sim (x \oplus y)$   
 $\langle proof \rangle$

**lemma** *xor-cancel-right* [*simp*]:  $x \oplus \sim x = \mathbf{1}$   
 $\langle proof \rangle$

**lemma** *xor-cancel-left* [*simp*]:  $\sim x \oplus x = \mathbf{1}$   
 $\langle proof \rangle$

**lemma** *conj-xor-distrib*:  $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$   
 $\langle proof \rangle$

**lemma** *conj-xor-distrib2*:  
 $(y \oplus z) \sqcap x = (y \sqcap x) \oplus (z \sqcap x)$   
 $\langle proof \rangle$

**lemmas** *conj-xor-distrib* =  
*conj-xor-distrib conj-xor-distrib2*

**end**

**end**

## 15 WordGenLib: Miscellaneous Library for Words

**theory** *WordGenLib*  
**imports** *WordShift Boolean-Algebra*  
**begin**

**declare** *of-nat-2p* [*simp*]

**lemma** *word-int-cases*:  
 $\llbracket \bigwedge n. \llbracket (x :: 'a :: \text{len } 0 \text{ word}) = \text{word-of-int } n; 0 \leq n; n < 2^{\text{len-of TYPE('a)}} \rrbracket \implies P \rrbracket$   
 $\implies P$   
 $\langle proof \rangle$

**lemma** *word-nat-cases* [*cases type: word*]:  
 $\llbracket \bigwedge n. \llbracket (x :: 'a :: \text{len word}) = \text{of-nat } n; n < 2^{\text{len-of TYPE('a)}} \rrbracket \implies P \rrbracket$   
 $\implies P$   
 $\langle proof \rangle$

**lemma** *max-word-eq*:  
 $(\text{max-word} :: 'a :: \text{len word}) = 2^{\text{len-of TYPE('a)}} - 1$   
 $\langle proof \rangle$

**lemma** *max-word-max* [*simp,intro!*]:

$$n \leq \text{max-word}$$

*<proof>*

**lemma** *word-of-int-2p-len*:

$$\text{word-of-int } (2 \wedge \text{len-of TYPE('a)}) = (0::'a::\text{len0 word})$$

*<proof>*

**lemma** *word-pow-0*:

$$(2::'a::\text{len word}) \wedge \text{len-of TYPE('a)} = 0$$

*<proof>*

**lemma** *max-word-wrap*:  $x + 1 = 0 \implies x = \text{max-word}$

*<proof>*

**lemma** *max-word-minus*:

$$\text{max-word} = (-1::'a::\text{len word})$$

*<proof>*

**lemma** *max-word-bl* [*simp*]:

$$\text{to-bl } (\text{max-word}::'a::\text{len word}) = \text{replicate } (\text{len-of TYPE('a)}) \text{ True}$$

*<proof>*

**lemma** *max-test-bit* [*simp*]:

$$(\text{max-word}::'a::\text{len word}) !! n = (n < \text{len-of TYPE('a)})$$

*<proof>*

**lemma** *word-and-max* [*simp*]:

$$x \text{ AND } \text{max-word} = x$$

*<proof>*

**lemma** *word-or-max* [*simp*]:

$$x \text{ OR } \text{max-word} = \text{max-word}$$

*<proof>*

**lemma** *word-ao-dist2*:

$$x \text{ AND } (y \text{ OR } z) = x \text{ AND } y \text{ OR } x \text{ AND } (z::'a::\text{len0 word})$$

*<proof>*

**lemma** *word-oa-dist2*:

$$x \text{ OR } y \text{ AND } z = (x \text{ OR } y) \text{ AND } (x \text{ OR } (z::'a::\text{len0 word}))$$

*<proof>*

**lemma** *word-and-not* [*simp*]:

$$x \text{ AND } \text{NOT } x = (0::'a::\text{len0 word})$$

*<proof>*

**lemma** *word-or-not* [*simp*]:

$$x \text{ OR } \text{NOT } x = \text{max-word}$$

$\langle \text{proof} \rangle$

**lemma** *word-boolean*:

*boolean* (*op AND*) (*op OR*) *bitNOT 0 max-word*  
 $\langle \text{proof} \rangle$

**interpretation** *word-bool-alg*:

*boolean* [*op AND op OR bitNOT 0 max-word*]  
 $\langle \text{proof} \rangle$

**lemma** *word-xor-and-or*:

$x \text{ XOR } y = x \text{ AND NOT } y \text{ OR NOT } x \text{ AND } (y::'a::\text{len0 word})$   
 $\langle \text{proof} \rangle$

**interpretation** *word-bool-alg*:

*boolean-xor* [*op AND op OR bitNOT 0 max-word op XOR*]  
 $\langle \text{proof} \rangle$

**lemma** *shiftr-0 [iff]*:

$(x::'a::\text{len0 word}) >> 0 = x$   
 $\langle \text{proof} \rangle$

**lemma** *shiftr-0 [simp]*:

$(x :: 'a :: \text{len word}) << 0 = x$   
 $\langle \text{proof} \rangle$

**lemma** *shiftr-1 [simp]*:

$(1::'a::\text{len word}) << n = 2^n$   
 $\langle \text{proof} \rangle$

**lemma** *uint-lt-0 [simp]*:

*uint*  $x < 0 = \text{False}$   
 $\langle \text{proof} \rangle$

**lemma** *shiftr1-1 [simp]*:

*shiftr1*  $(1::'a::\text{len word}) = 0$   
 $\langle \text{proof} \rangle$

**lemma** *shiftr-1 [simp]*:

$(1::'a::\text{len word}) >> n = (\text{if } n = 0 \text{ then } 1 \text{ else } 0)$   
 $\langle \text{proof} \rangle$

**lemma** *word-less-1 [simp]*:

$((x::'a::\text{len word}) < 1) = (x = 0)$   
 $\langle \text{proof} \rangle$

**lemma** *to-bl-mask*:

*to-bl*  $(\text{mask } n :: 'a::\text{len word}) =$   
 $\text{replicate } (\text{len-of TYPE('a)} - n) \text{ False } @$

*replicate (min (len-of TYPE('a)) n) True*  
 ⟨proof⟩

**lemma** *map-replicate-True*:  
*n = length xs ==>*  
*map (λ(x,y). x & y) (zip xs (replicate n True)) = xs*  
 ⟨proof⟩

**lemma** *map-replicate-False*:  
*n = length xs ==> map (λ(x,y). x & y)*  
*(zip xs (replicate n False)) = replicate n False*  
 ⟨proof⟩

**lemma** *bl-and-mask*:  
**fixes** *w :: 'a::len word*  
**fixes** *n*  
**defines** *n' ≡ len-of TYPE('a) - n*  
**shows** *to-bl (w AND mask n) = replicate n' False @ drop n' (to-bl w)*  
 ⟨proof⟩

**lemma** *drop-rev-takefill*:  
*length xs ≤ n ==>*  
*drop (n - length xs) (rev (takefill False n (rev xs))) = xs*  
 ⟨proof⟩

**lemma** *map-nth-0* [simp]:  
*map (op !! (0::'a::len0 word)) xs = replicate (length xs) False*  
 ⟨proof⟩

**lemma** *uint-plus-if-size*:  
*uint (x + y) =*  
*(if uint x + uint y < 2^size x then*  
*uint x + uint y*  
*else*  
*uint x + uint y - 2^size x)*  
 ⟨proof⟩

**lemma** *unat-plus-if-size*:  
*unat (x + (y::'a::len word)) =*  
*(if unat x + unat y < 2^size x then*  
*unat x + unat y*  
*else*  
*unat x + unat y - 2^size x)*  
 ⟨proof⟩

**lemma** *word-neq-0-conv* [simp]:  
**fixes** *w :: 'a :: len word*  
**shows** *(w ≠ 0) = (0 < w)*  
 ⟨proof⟩

**lemma** *max-lt*:

$unat (max\ a\ b\ div\ c) = unat (max\ a\ b)\ div\ unat\ (c :: 'a :: len\ word)$   
 $\langle proof \rangle$

**lemma** *uint-sub-if-size*:

$uint\ (x - y) =$   
 $(if\ uint\ y \leq uint\ x\ then$   
 $\quad uint\ x - uint\ y$   
 $else$   
 $\quad uint\ x - uint\ y + 2^{size\ x})$   
 $\langle proof \rangle$

**lemma** *unat-sub-simple*:

$x \leq y \implies unat\ (y - x) = unat\ y - unat\ x$   
 $\langle proof \rangle$

**lemmas** *unat-sub = unat-sub-simple*

**lemma** *word-less-sub1*:

**fixes**  $x :: 'a :: len\ word$   
**shows**  $x \neq 0 \implies 1 < x = (0 < x - 1)$   
 $\langle proof \rangle$

**lemma** *word-le-sub1*:

**fixes**  $x :: 'a :: len\ word$   
**shows**  $x \neq 0 \implies 1 \leq x = (0 \leq x - 1)$   
 $\langle proof \rangle$

**lemmas** *word-less-sub1-numberof [simp] =*  
*word-less-sub1 [of number-of w, standard]*

**lemmas** *word-le-sub1-numberof [simp] =*  
*word-le-sub1 [of number-of w, standard]*

**lemma** *word-of-int-minus*:

$word-of-int\ (2^{len-of\ TYPE('a)} - i) = (word-of-int\ (-i) :: 'a :: len\ word)$   
 $\langle proof \rangle$

**lemmas** *word-of-int-inj =*

*word-uint.Abs-inject [unfolded uints-num, simplified]*

**lemma** *word-le-less-eq*:

$(x :: 'z :: len\ word) \leq y = (x = y \vee x < y)$   
 $\langle proof \rangle$

**lemma** *mod-plus-cong*:

**assumes**  $1: (b :: int) = b'$   
**and**  $2: x \bmod b' = x' \bmod b'$   
**and**  $3: y \bmod b' = y' \bmod b'$

**and**  $4: x' + y' = z'$   
**shows**  $(x + y) \bmod b = z' \bmod b'$   
 $\langle \text{proof} \rangle$

**lemma** *mod-minus-cong*:  
**assumes**  $1: (b::\text{int}) = b'$   
**and**  $2: x \bmod b' = x' \bmod b'$   
**and**  $3: y \bmod b' = y' \bmod b'$   
**and**  $4: x' - y' = z'$   
**shows**  $(x - y) \bmod b = z' \bmod b'$   
 $\langle \text{proof} \rangle$

**lemma** *word-induct-less*:  
 $\llbracket P (0::'a::\text{len word}); \bigwedge n. \llbracket n < m; P n \rrbracket \implies P (1 + n) \rrbracket \implies P m$   
 $\langle \text{proof} \rangle$

**lemma** *word-induct*:  
 $\llbracket P (0::'a::\text{len word}); \bigwedge n. P n \implies P (1 + n) \rrbracket \implies P m$   
 $\langle \text{proof} \rangle$

**lemma** *word-induct2* [*induct type*]:  
 $\llbracket P 0; \bigwedge n. \llbracket 1 + n \neq 0; P n \rrbracket \implies P (1 + n) \rrbracket \implies P (n::'b::\text{len word})$   
 $\langle \text{proof} \rangle$

**constdefs**  
 $\text{word-rec} :: 'a \Rightarrow ('b::\text{len word} \Rightarrow 'a \Rightarrow 'a) \Rightarrow 'b \text{ word} \Rightarrow 'a$   
 $\text{word-rec forZero forSuc } n \equiv \text{nat-rec forZero (forSuc } \circ \text{ of-nat)} (\text{unat } n)$

**lemma** *word-rec-0*:  $\text{word-rec } z \ s \ 0 = z$   
 $\langle \text{proof} \rangle$

**lemma** *word-rec-Suc*:  
 $1 + n \neq (0::'a::\text{len word}) \implies \text{word-rec } z \ s \ (1 + n) = s \ n \ (\text{word-rec } z \ s \ n)$   
 $\langle \text{proof} \rangle$

**lemma** *word-rec-Pred*:  
 $n \neq 0 \implies \text{word-rec } z \ s \ n = s \ (n - 1) \ (\text{word-rec } z \ s \ (n - 1))$   
 $\langle \text{proof} \rangle$

**lemma** *word-rec-in*:  
 $f \ (\text{word-rec } z \ (\lambda-. f) \ n) = \text{word-rec } (f \ z) \ (\lambda-. f) \ n$   
 $\langle \text{proof} \rangle$

**lemma** *word-rec-in2*:  
 $f \ n \ (\text{word-rec } z \ f \ n) = \text{word-rec } (f \ 0 \ z) \ (f \circ \text{op} + 1) \ n$   
 $\langle \text{proof} \rangle$

**lemma** *word-rec-twice*:  
 $m \leq n \implies \text{word-rec } z \ f \ n = \text{word-rec } (\text{word-rec } z \ f \ (n - m)) \ (f \circ \text{op} + (n -$

$m))\ m$   
 $\langle proof \rangle$

**lemma** *word-rec-id*: *word-rec*  $z\ (\lambda\ -. id)\ n = z$   
 $\langle proof \rangle$

**lemma** *word-rec-id-eq*:  $\forall m < n. f\ m = id \implies word-rec\ z\ f\ n = z$   
 $\langle proof \rangle$

**lemma** *word-rec-max*:  
 $\forall m \geq n. m \neq -1 \implies f\ m = id \implies word-rec\ z\ f\ -1 = word-rec\ z\ f\ n$   
 $\langle proof \rangle$

**lemma** *unatSuc*:  
 $1 + n \neq (0::'a::len\ word) \implies unat\ (1 + n) = Suc\ (unat\ n)$   
 $\langle proof \rangle$

**end**

## 16 WordMain: Main Word Library

**theory** *WordMain*  
**imports** *WordGenLib*  
**begin**

**lemmas** *word-no-1* [*simp*] = *word-1-no* [*symmetric, unfolded BIT-simps*]  
**lemmas** *word-no-0* [*simp*] = *word-0-no* [*symmetric*]

**declare** *word-0-bl* [*simp*]  
**declare** *bin-to-bl-def* [*simp*]  
**declare** *to-bl-0* [*simp*]  
**declare** *of-bl-True* [*simp*]

Examples

**types** *word32* = 32 *word*  
**types** *word8* = 8 *word*  
**types** *byte* = *word8*

for more see *WordExamples.thy*

**end**

## References

- [1] Jeremy Dawson. Isabelle theories for machine words. In Michael Goldsmith and Bill Roscoe, editors, *Seventh International Workshop on Au-*



*tomated Verification of Critical Systems (AVOCS'07)*, Electronic Notes in Theoretical Computer Science, page 15, Oxford, September 2007. Elsevier. to appear.