

# Matrix

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```
theory MatrixGeneral
imports Main
begin
```

```
types 'a infmatrix = [nat, nat]  $\Rightarrow$  'a
```

```
constdefs
```

```
  nonzero-positions :: ('a::zero) infmatrix  $\Rightarrow$  (nat*nat) set
  nonzero-positions A == {pos. A (fst pos) (snd pos)  $\sim$  0}
```

```
typedef 'a matrix = {(f::('a::zero) infmatrix)). finite (nonzero-positions f)}
  <proof>
```

```
declare Rep-matrix-inverse[simp]
```

```
lemma finite-nonzero-positions : finite (nonzero-positions (Rep-matrix A))
  <proof>
```

```
constdefs
```

```
  nrows :: ('a::zero) matrix  $\Rightarrow$  nat
  nrows A == if nonzero-positions(Rep-matrix A) = {} then 0 else Suc(Max
    ((image fst) (nonzero-positions (Rep-matrix A))))
  ncols :: ('a::zero) matrix  $\Rightarrow$  nat
  ncols A == if nonzero-positions(Rep-matrix A) = {} then 0 else Suc(Max ((image
    snd) (nonzero-positions (Rep-matrix A))))
```

```
lemma nrows:
```

```
  assumes hyp: nrows A  $\leq$  m
  shows (Rep-matrix A m n) = 0 (is ?concl)
  <proof>
```

```
constdefs
```

```
  transpose-infmatrix :: 'a infmatrix  $\Rightarrow$  'a infmatrix
  transpose-infmatrix A j i == A i j
  transpose-matrix :: ('a::zero) matrix  $\Rightarrow$  'a matrix
  transpose-matrix == Abs-matrix o transpose-infmatrix o Rep-matrix
```

**declare** *transpose-infmatrix-def*[simp]

**lemma** *transpose-infmatrix-twice*[simp]: *transpose-infmatrix* (*transpose-infmatrix* *A*) = *A*  
 <proof>

**lemma** *transpose-infmatrix*: *transpose-infmatrix* (% *j i*. *P j i*) = (% *j i*. *P i j*)  
 <proof>

**lemma** *transpose-infmatrix-closed*[simp]: *Rep-matrix* (*Abs-matrix* (*transpose-infmatrix* (*Rep-matrix* *x*))) = *transpose-infmatrix* (*Rep-matrix* *x*)  
 <proof>

**lemma** *infmatrixforward*: (*x::'a infmatrix*) = *y*  $\implies \forall a b. x a b = y a b$  <proof>

**lemma** *transpose-infmatrix-inject*: (*transpose-infmatrix* *A* = *transpose-infmatrix* *B*) = (*A* = *B*)  
 <proof>

**lemma** *transpose-matrix-inject*: (*transpose-matrix* *A* = *transpose-matrix* *B*) = (*A* = *B*)  
 <proof>

**lemma** *transpose-matrix*[simp]: *Rep-matrix*(*transpose-matrix* *A*) *j i* = *Rep-matrix* *A i j*  
 <proof>

**lemma** *transpose-transpose-id*[simp]: *transpose-matrix* (*transpose-matrix* *A*) = *A*  
 <proof>

**lemma** *nrows-transpose*[simp]: *nrows* (*transpose-matrix* *A*) = *ncols* *A*  
 <proof>

**lemma** *ncols-transpose*[simp]: *ncols* (*transpose-matrix* *A*) = *nrows* *A*  
 <proof>

**lemma** *ncols*: *ncols* *A* <= *n*  $\implies$  *Rep-matrix* *A m n* = 0  
 <proof>

**lemma** *ncols-le*: (*ncols* *A* <= *n*) = (! *j i*. *n* <= *i*  $\longrightarrow$  (*Rep-matrix* *A j i*) = 0) (is  
 - = ?st)  
 <proof>

**lemma** *less-ncols*: (*n* < *ncols* *A*) = (? *j i*. *n* <= *i* & (*Rep-matrix* *A j i*)  $\neq$  0) (is  
 ?concl)  
 <proof>

**lemma** *le-ncols*: (*n* <= *ncols* *A*) = ( $\forall m. (\forall j i. m \leq i \longrightarrow (\text{Rep-matrix } A j i))$ )

$= 0) \longrightarrow n \leq m$  (**is** ?concl)  
 <proof>

**lemma** *nrows-le*:  $(\text{nrows } A \leq n) = (! j i. n \leq j \longrightarrow (\text{Rep-matrix } A j i) = 0)$   
 (**is** ?s)  
 <proof>

**lemma** *less-nrows*:  $(m < \text{nrows } A) = (? j i. m \leq j \ \& \ (\text{Rep-matrix } A j i) \neq 0)$   
 (**is** ?concl)  
 <proof>

**lemma** *le-nrows*:  $(n \leq \text{nrows } A) = (\forall m. (\forall j i. m \leq j \longrightarrow (\text{Rep-matrix } A j i) = 0) \longrightarrow n \leq m)$  (**is** ?concl)  
 <proof>

**lemma** *nrows-notzero*:  $\text{Rep-matrix } A m n \neq 0 \implies m < \text{nrows } A$   
 <proof>

**lemma** *ncols-notzero*:  $\text{Rep-matrix } A m n \neq 0 \implies n < \text{ncols } A$   
 <proof>

**lemma** *finite-natarray1*:  $\text{finite } \{x. x < (n::\text{nat})\}$   
 <proof>

**lemma** *finite-natarray2*:  $\text{finite } \{\text{pos}. (\text{fst pos}) < (m::\text{nat}) \ \& \ (\text{snd pos}) < (n::\text{nat})\}$   
 <proof>

**lemma** *RepAbs-matrix*:  
**assumes** *aem*:  $? m. ! j i. m \leq j \longrightarrow x j i = 0$  (**is** ?em) **and** *aen*:  $? n. ! j i. (n \leq i \longrightarrow x j i = 0)$  (**is** ?en)  
**shows**  $(\text{Rep-matrix } (\text{Abs-matrix } x)) = x$   
 <proof>

**constdefs**

*apply-infmatrix* ::  $('a \Rightarrow 'b) \Rightarrow 'a \text{ infmatrix} \Rightarrow 'b \text{ infmatrix}$   
*apply-infmatrix* f == % A. (% j i. f (A j i))  
*apply-matrix* ::  $('a \Rightarrow 'b) \Rightarrow ('a::\text{zero}) \text{ matrix} \Rightarrow ('b::\text{zero}) \text{ matrix}$   
*apply-matrix* f == % A. Abs-matrix (apply-infmatrix f (Rep-matrix A))  
*combine-infmatrix* ::  $('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow 'a \text{ infmatrix} \Rightarrow 'b \text{ infmatrix} \Rightarrow 'c \text{ infmatrix}$   
*combine-infmatrix* f == % A B. (% j i. f (A j i) (B j i))  
*combine-matrix* ::  $('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow ('a::\text{zero}) \text{ matrix} \Rightarrow ('b::\text{zero}) \text{ matrix} \Rightarrow ('c::\text{zero}) \text{ matrix}$   
*combine-matrix* f == % A B. Abs-matrix (combine-infmatrix f (Rep-matrix A) (Rep-matrix B))

**lemma** *expand-apply-infmatrix[simp]*:  $\text{apply-infmatrix } f A j i = f (A j i)$   
 <proof>

**lemma** *expand-combine-infmatrix[simp]*:  $\text{combine-infmatrix } f A B j i = f (A j i)$

(B j i)  
 <proof>

**constdefs**

*commutative* :: ('a  $\Rightarrow$  'a  $\Rightarrow$  'b)  $\Rightarrow$  bool  
*commutative* f == ! x y. f x y = f y x  
*associative* :: ('a  $\Rightarrow$  'a  $\Rightarrow$  'a)  $\Rightarrow$  bool  
*associative* f == ! x y z. f (f x y) z = f x (f y z)

To reason about associativity and commutativity of operations on matrices, let's take a step back and look at the general situation: Assume that we have sets  $A$  and  $B$  with  $B \subset A$  and an abstraction  $u : A \rightarrow B$ . This abstraction has to fulfill  $u(b) = b$  for all  $b \in B$ , but is arbitrary otherwise. Each function  $f : A \times A \rightarrow A$  now induces a function  $f' : B \times B \rightarrow B$  by  $f' = u \circ f$ . It is obvious that commutativity of  $f$  implies commutativity of  $f'$ :  $f'xy = u(fxy) = u(fyx) = f'yx$ .

**lemma** *combine-infmatrix-commute*:

*commutative* f  $\implies$  *commutative* (combine-infmatrix f)  
 <proof>

**lemma** *combine-matrix-commute*:

*commutative* f  $\implies$  *commutative* (combine-matrix f)  
 <proof>

On the contrary, given an associative function  $f$  we cannot expect  $f'$  to be associative. A counterexample is given by  $A = \mathbb{Z}$ ,  $B = \{-1, 0, 1\}$ , as  $f$  we take addition on  $\mathbb{Z}$ , which is clearly associative. The abstraction is given by  $u(a) = 0$  for  $a \notin B$ . Then we have

$$f'(f'11) - 1 = u(f(u(f11)) - 1) = u(f(u2) - 1) = u(f0 - 1) = -1,$$

but on the other hand we have

$$f'1(f'1 - 1) = u(f1(u(f1 - 1))) = u(f10) = 1.$$

A way out of this problem is to assume that  $f(A \times A) \subset A$  holds, and this is what we are going to do:

**lemma** *nonzero-positions-combine-infmatrix[simp]*:  $f \ 0 \ 0 = 0 \implies \text{nonzero-positions } (\text{combine-infmatrix } f \ A \ B) \subseteq (\text{nonzero-positions } A) \cup (\text{nonzero-positions } B)$   
 <proof>

**lemma** *finite-nonzero-positions-Rep[simp]*: *finite* (nonzero-positions (Rep-matrix A))  
 <proof>

**lemma** *combine-infmatrix-closed [simp]*:

$f \ 0 \ 0 = 0 \implies \text{Rep-matrix } (\text{Abs-matrix } (\text{combine-infmatrix } f \ (\text{Rep-matrix } A) \ (\text{Rep-matrix } B))) = \text{combine-infmatrix } f \ (\text{Rep-matrix } A) \ (\text{Rep-matrix } B)$

*<proof>*

We need the next two lemmas only later, but it is analog to the above one, so we prove them now:

**lemma** *nonzero-positions-apply-infmatrix[simp]:*  $f \ 0 = 0 \implies \text{nonzero-positions } (\text{apply-infmatrix } f \ A) \subseteq \text{nonzero-positions } A$   
*<proof>*

**lemma** *apply-infmatrix-closed [simp]:*  
 $f \ 0 = 0 \implies \text{Rep-matrix } (\text{Abs-matrix } (\text{apply-infmatrix } f \ (\text{Rep-matrix } A))) = \text{apply-infmatrix } f \ (\text{Rep-matrix } A)$   
*<proof>*

**lemma** *combine-infmatrix-assoc[simp]:*  $f \ 0 \ 0 = 0 \implies \text{associative } f \implies \text{associative } (\text{combine-infmatrix } f)$   
*<proof>*

**lemma** *comb:*  $f = g \implies x = y \implies f \ x = g \ y$   
*<proof>*

**lemma** *combine-matrix-assoc:*  $f \ 0 \ 0 = 0 \implies \text{associative } f \implies \text{associative } (\text{combine-matrix } f)$   
*<proof>*

**lemma** *Rep-apply-matrix[simp]:*  $f \ 0 = 0 \implies \text{Rep-matrix } (\text{apply-matrix } f \ A) \ j \ i = f \ (\text{Rep-matrix } A \ j \ i)$   
*<proof>*

**lemma** *Rep-combine-matrix[simp]:*  $f \ 0 \ 0 = 0 \implies \text{Rep-matrix } (\text{combine-matrix } f \ A \ B) \ j \ i = f \ (\text{Rep-matrix } A \ j \ i) \ (\text{Rep-matrix } B \ j \ i)$   
*<proof>*

**lemma** *combine-nrows-max:*  $f \ 0 \ 0 = 0 \implies \text{nrows } (\text{combine-matrix } f \ A \ B) \leq \max (\text{nrows } A) (\text{nrows } B)$   
*<proof>*

**lemma** *combine-ncols-max:*  $f \ 0 \ 0 = 0 \implies \text{ncols } (\text{combine-matrix } f \ A \ B) \leq \max (\text{ncols } A) (\text{ncols } B)$   
*<proof>*

**lemma** *combine-nrows:*  $f \ 0 \ 0 = 0 \implies \text{nrows } A \leq q \implies \text{nrows } B \leq q \implies \text{nrows } (\text{combine-matrix } f \ A \ B) \leq q$   
*<proof>*

**lemma** *combine-ncols:*  $f \ 0 \ 0 = 0 \implies \text{ncols } A \leq q \implies \text{ncols } B \leq q \implies \text{ncols } (\text{combine-matrix } f \ A \ B) \leq q$   
*<proof>*

**constdefs**

$zero\text{-}r\text{-}neutral :: ('a \Rightarrow 'b :: zero \Rightarrow 'a) \Rightarrow bool$   
 $zero\text{-}r\text{-}neutral\ f == !\ a.\ f\ a\ 0 = a$   
 $zero\text{-}l\text{-}neutral :: ('a :: zero \Rightarrow 'b \Rightarrow 'a) \Rightarrow bool$   
 $zero\text{-}l\text{-}neutral\ f == !\ a.\ f\ 0\ a = a$   
 $zero\text{-}closed :: (('a :: zero) \Rightarrow ('b :: zero) \Rightarrow ('c :: zero)) \Rightarrow bool$   
 $zero\text{-}closed\ f == (!x.\ f\ x\ 0 = 0) \ \&\ (\!y.\ f\ 0\ y = 0)$

**consts**  $foldseq :: ('a \Rightarrow 'a \Rightarrow 'a) \Rightarrow (nat \Rightarrow 'a) \Rightarrow nat \Rightarrow 'a$   
**primrec**  
 $foldseq\ f\ s\ 0 = s\ 0$   
 $foldseq\ f\ s\ (Suc\ n) = f\ (s\ 0)\ (foldseq\ f\ (\% k.\ s(Suc\ k))\ n)$

**consts**  $foldseq\text{-}transposed :: ('a \Rightarrow 'a \Rightarrow 'a) \Rightarrow (nat \Rightarrow 'a) \Rightarrow nat \Rightarrow 'a$   
**primrec**  
 $foldseq\text{-}transposed\ f\ s\ 0 = s\ 0$   
 $foldseq\text{-}transposed\ f\ s\ (Suc\ n) = f\ (foldseq\text{-}transposed\ f\ s\ n)\ (s\ (Suc\ n))$

**lemma**  $foldseq\text{-}assoc : associative\ f \Longrightarrow foldseq\ f = foldseq\text{-}transposed\ f$   
 $\langle proof \rangle$

**lemma**  $foldseq\text{-}distr : \llbracket associative\ f ; commutative\ f \rrbracket \Longrightarrow foldseq\ f\ (\% k.\ f\ (u\ k)\ (v\ k))\ n = f\ (foldseq\ f\ u\ n)\ (foldseq\ f\ v\ n)$   
 $\langle proof \rangle$

**theorem**  $\llbracket associative\ f ; associative\ g ; \forall a\ b\ c\ d.\ g\ (f\ a\ b)\ (f\ c\ d) = f\ (g\ a\ c)\ (g\ b\ d) ; ?\ x\ y.\ (f\ x) \neq (f\ y) ; ?\ x\ y.\ (g\ x) \neq (g\ y) ; f\ x\ x = x ; g\ x\ x = x \rrbracket \Longrightarrow f=g \mid (!\ y.\ f\ y\ y = y) \mid (!\ y.\ g\ y\ y = y)$   
 $\langle proof \rangle$

**lemma**  $foldseq\text{-}zero :$   
**assumes**  $fz : f\ 0\ 0 = 0$  **and**  $sz : !\ i.\ i \leq n \longrightarrow s\ i = 0$   
**shows**  $foldseq\ f\ s\ n = 0$   
 $\langle proof \rangle$

**lemma**  $foldseq\text{-}significant\text{-}positions :$   
**assumes**  $p : !\ i.\ i \leq N \longrightarrow S\ i = T\ i$   
**shows**  $foldseq\ f\ S\ N = foldseq\ f\ T\ N$  (**is**  $?concl$ )  
 $\langle proof \rangle$

**lemma**  $foldseq\text{-}tail : M \leq N \Longrightarrow foldseq\ f\ S\ N = foldseq\ f\ (\% k.\ (if\ k < M\ then\ (S\ k)\ else\ (foldseq\ f\ (\% k.\ S(k+M))\ (N-M))))\ M$  (**is**  $?p \Longrightarrow ?concl$ )  
 $\langle proof \rangle$

**lemma**  $foldseq\text{-}zerotail :$   
**assumes**  
 $fz : f\ 0\ 0 = 0$   
**and**  $sz : !\ i.\ n \leq i \longrightarrow s\ i = 0$   
**and**  $nm : n \leq m$

**shows**  
 $foldseq\ f\ s\ n = foldseq\ f\ s\ m$   
 $\langle proof \rangle$

**lemma** *foldseq-zerotail2*:  
**assumes**  $! x. f\ x\ 0 = x$   
**and**  $! i. n < i \longrightarrow s\ i = 0$   
**and**  $nm: n \leq m$   
**shows**  
 $foldseq\ f\ s\ n = foldseq\ f\ s\ m$  (**is**  $?concl$ )  
 $\langle proof \rangle$

**lemma** *foldseq-zerostart*:  
 $! x. f\ 0\ (f\ 0\ x) = f\ 0\ x \implies ! i. i \leq n \longrightarrow s\ i = 0 \implies foldseq\ f\ s\ (Suc\ n) = f\ 0\ (s\ (Suc\ n))$   
 $\langle proof \rangle$

**lemma** *foldseq-zerostart2*:  
 $! x. f\ 0\ x = x \implies ! i. i < n \longrightarrow s\ i = 0 \implies foldseq\ f\ s\ n = s\ n$   
 $\langle proof \rangle$

**lemma** *foldseq-almostzero*:  
**assumes**  $f0x: ! x. f\ 0\ x = x$  **and**  $fx0: ! x. f\ x\ 0 = x$  **and**  $s0: ! i. i \neq j \longrightarrow s\ i = 0$   
**shows**  $foldseq\ f\ s\ n = (if\ (j \leq n)\ then\ (s\ j)\ else\ 0)$  (**is**  $?concl$ )  
 $\langle proof \rangle$

**lemma** *foldseq-distr-unary*:  
**assumes**  $!! a\ b. g\ (f\ a\ b) = f\ (g\ a)\ (g\ b)$   
**shows**  $g(foldseq\ f\ s\ n) = foldseq\ f\ (\% x. g(s\ x))\ n$  (**is**  $?concl$ )  
 $\langle proof \rangle$

**constdefs**  
 $mult\_matrix\_n :: nat \Rightarrow ((a::zero) \Rightarrow (b::zero) \Rightarrow (c::zero)) \Rightarrow ('c \Rightarrow 'c \Rightarrow 'c)$   
 $\Rightarrow 'a\ matrix \Rightarrow 'b\ matrix \Rightarrow 'c\ matrix$   
 $mult\_matrix\_n\ n\ fmul\ fadd\ A\ B == Abs\_matrix(\% j\ i. foldseq\ fadd\ (\% k. fmul\ (Rep\_matrix\ A\ j\ k)\ (Rep\_matrix\ B\ k\ i))\ n)$   
 $mult\_matrix :: ((a::zero) \Rightarrow (b::zero) \Rightarrow (c::zero)) \Rightarrow ('c \Rightarrow 'c \Rightarrow 'c) \Rightarrow 'a\ matrix \Rightarrow 'b\ matrix \Rightarrow 'c\ matrix$   
 $mult\_matrix\ fmul\ fadd\ A\ B == mult\_matrix\_n\ (max\ (ncols\ A)\ (nrows\ B))\ fmul\ fadd\ A\ B$

**lemma** *mult-matrix-n*:  
**assumes**  $prems: ncols\ A \leq n$  (**is**  $?An$ )  $nrows\ B \leq n$  (**is**  $?Bn$ )  $fadd\ 0\ 0 = 0$   $fmul\ 0\ 0 = 0$   
**shows**  $c: mult\_matrix\ fmul\ fadd\ A\ B = mult\_matrix\_n\ n\ fmul\ fadd\ A\ B$  (**is**  $?concl$ )  
 $\langle proof \rangle$

**lemma** *mult-matrix-nm*:  
**assumes**  $prems: ncols\ A \leq n$   $nrows\ B \leq n$   $ncols\ A \leq m$   $nrows\ B \leq m$

*fadd 0 0 = 0 fmul 0 0 = 0*  
**shows** *mult-matrix-n n fmul fadd A B = mult-matrix-n m fmul fadd A B*  
*<proof>*

**constdefs**

*r-distributive* :: (*'a* ⇒ *'b* ⇒ *'b*) ⇒ (*'b* ⇒ *'b* ⇒ *'b*) ⇒ *bool*  
*r-distributive fmul fadd* == ! *a u v. fmul a (fadd u v) = fadd (fmul a u) (fmul a v)*  
*l-distributive* :: (*'a* ⇒ *'b* ⇒ *'a*) ⇒ (*'a* ⇒ *'a* ⇒ *'a*) ⇒ *bool*  
*l-distributive fmul fadd* == ! *a u v. fmul (fadd u v) a = fadd (fmul u a) (fmul v a)*  
*distributive* :: (*'a* ⇒ *'a* ⇒ *'a*) ⇒ (*'a* ⇒ *'a* ⇒ *'a*) ⇒ *bool*  
*distributive fmul fadd* == *l-distributive fmul fadd & r-distributive fmul fadd*

**lemma** *max1*: !! *a x y. (a::nat) <= x ⇒ a <= max x y* *<proof>*

**lemma** *max2*: !! *b x y. (b::nat) <= y ⇒ b <= max x y* *<proof>*

**lemma** *r-distributive-matrix*:

**assumes** *prems*:  
*r-distributive fmul fadd*  
*associative fadd*  
*commutative fadd*  
*fadd 0 0 = 0*  
! *a. fmul a 0 = 0*  
! *a. fmul 0 a = 0*  
**shows** *r-distributive (mult-matrix fmul fadd) (combine-matrix fadd)* (**is** *?concl*)  
*<proof>*

**lemma** *l-distributive-matrix*:

**assumes** *prems*:  
*l-distributive fmul fadd*  
*associative fadd*  
*commutative fadd*  
*fadd 0 0 = 0*  
! *a. fmul a 0 = 0*  
! *a. fmul 0 a = 0*  
**shows** *l-distributive (mult-matrix fmul fadd) (combine-matrix fadd)* (**is** *?concl*)  
*<proof>*

**instance** *matrix* :: (*zero*) *zero* *<proof>*

**defs**(**overloaded**)

*zero-matrix-def*: (*0::('a::zero) matrix*) == *Abs-matrix(% j i. 0)*

**lemma** *Rep-zero-matrix-def[simp]*: *Rep-matrix 0 j i = 0*  
*<proof>*

**lemma** *zero-matrix-def-nrows[simp]*: *nrows 0 = 0*  
*<proof>*



**lemma** *zero-matrix-def-ncols*[simp]:  $\text{ncols } 0 = 0$

*<proof>*

**lemma** *combine-matrix-zero-l-neutral*:  $\text{zero-l-neutral } f \implies \text{zero-l-neutral } (\text{combine-matrix } f)$

*<proof>*

**lemma** *combine-matrix-zero-r-neutral*:  $\text{zero-r-neutral } f \implies \text{zero-r-neutral } (\text{combine-matrix } f)$

*<proof>*

**lemma** *mult-matrix-zero-closed*:  $\llbracket \text{fadd } 0 \ 0 = 0; \text{zero-closed } \text{fmul} \rrbracket \implies \text{zero-closed}$   
(*mult-matrix fmul fadd*)

*<proof>*

**lemma** *mult-matrix-n-zero-right*[simp]:  $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } a \ 0 = 0 \rrbracket \implies$   
*mult-matrix-n n fmul fadd A 0 = 0*

*<proof>*

**lemma** *mult-matrix-n-zero-left*[simp]:  $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } 0 \ a = 0 \rrbracket \implies$   
*mult-matrix-n n fmul fadd 0 A = 0*

*<proof>*

**lemma** *mult-matrix-zero-left*[simp]:  $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } 0 \ a = 0 \rrbracket \implies \text{mult-matrix}$   
*fmul fadd 0 A = 0*

*<proof>*

**lemma** *mult-matrix-zero-right*[simp]:  $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } a \ 0 = 0 \rrbracket \implies \text{mult-matrix}$   
*fmul fadd A 0 = 0*

*<proof>*

**lemma** *apply-matrix-zero*[simp]:  $f \ 0 = 0 \implies \text{apply-matrix } f \ 0 = 0$

*<proof>*

**lemma** *combine-matrix-zero*:  $f \ 0 \ 0 = 0 \implies \text{combine-matrix } f \ 0 \ 0 = 0$

*<proof>*

**lemma** *transpose-matrix-zero*[simp]:  $\text{transpose-matrix } 0 = 0$

*<proof>*

**lemma** *apply-zero-matrix-def*[simp]:  $\text{apply-matrix } (\% \ x. \ 0) \ A = 0$

*<proof>*

**constdefs**

*singleton-matrix* ::  $\text{nat} \Rightarrow \text{nat} \Rightarrow ('a::\text{zero}) \Rightarrow 'a \ \text{matrix}$

*singleton-matrix*  $j \ i \ a == \text{Abs-matrix}(\% \ m \ n. \text{if } j = m \ \& \ i = n \text{ then } a \text{ else } 0)$

*move-matrix* ::  $('a::\text{zero}) \ \text{matrix} \Rightarrow \text{int} \Rightarrow \text{int} \Rightarrow 'a \ \text{matrix}$

*move-matrix*  $A \ y \ x == \text{Abs-matrix}(\% \ j \ i. \text{if } (\text{neg } ((\text{int } j) - y)) \mid (\text{neg } ((\text{int } i) - x)))$

then 0 else Rep-matrix A (nat ((int j)−y)) (nat ((int i)−x)))  
 take-rows :: ('a::zero) matrix ⇒ nat ⇒ 'a matrix  
 take-rows A r == Abs-matrix(% j i. if (j < r) then (Rep-matrix A j i) else 0)  
 take-columns :: ('a::zero) matrix ⇒ nat ⇒ 'a matrix  
 take-columns A c == Abs-matrix(% j i. if (i < c) then (Rep-matrix A j i) else 0)

#### constdefs

column-of-matrix :: ('a::zero) matrix ⇒ nat ⇒ 'a matrix  
 column-of-matrix A n == take-columns (move-matrix A 0 (− int n)) 1  
 row-of-matrix :: ('a::zero) matrix ⇒ nat ⇒ 'a matrix  
 row-of-matrix A m == take-rows (move-matrix A (− int m) 0) 1

**lemma** Rep-singleton-matrix[simp]: Rep-matrix (singleton-matrix j i e) m n = (if j = m & i = n then e else 0)  
 <proof>

**lemma** apply-singleton-matrix[simp]: f 0 = 0 ⇒ apply-matrix f (singleton-matrix j i x) = (singleton-matrix j i (f x))  
 <proof>

**lemma** singleton-matrix-zero[simp]: singleton-matrix j i 0 = 0  
 <proof>

**lemma** nrows-singleton[simp]: nrows(singleton-matrix j i e) = (if e = 0 then 0 else Suc j)  
 <proof>

**lemma** ncols-singleton[simp]: ncols(singleton-matrix j i e) = (if e = 0 then 0 else Suc i)  
 <proof>

**lemma** combine-singleton: f 0 0 = 0 ⇒ combine-matrix f (singleton-matrix j i a) (singleton-matrix j i b) = singleton-matrix j i (f a b)  
 <proof>

**lemma** transpose-singleton[simp]: transpose-matrix (singleton-matrix j i a) = singleton-matrix i j a  
 <proof>

**lemma** Rep-move-matrix[simp]:  
 Rep-matrix (move-matrix A y x) j i =  
 (if (neg ((int j)−y)) | (neg ((int i)−x)) then 0 else Rep-matrix A (nat((int j)−y)) (nat((int i)−x)))  
 <proof>

**lemma** move-matrix-0-0[simp]: move-matrix A 0 0 = A  
 <proof>

**lemma** *move-matrix-ortho*:  $\text{move-matrix } A \ j \ i = \text{move-matrix } (\text{move-matrix } A \ j \ 0) \ 0 \ i$   
 <proof>

**lemma** *transpose-move-matrix[simp]*:  
 $\text{transpose-matrix } (\text{move-matrix } A \ x \ y) = \text{move-matrix } (\text{transpose-matrix } A) \ y \ x$   
 <proof>

**lemma** *move-matrix-singleton[simp]*:  $\text{move-matrix } (\text{singleton-matrix } u \ v \ x) \ j \ i =$   
 $(\text{if } (j + \text{int } u < 0) \mid (i + \text{int } v < 0) \text{ then } 0 \text{ else } (\text{singleton-matrix } (\text{nat } (j + \text{int } u)) \ (\text{nat } (i + \text{int } v)) \ x))$   
 <proof>

**lemma** *Rep-take-columns[simp]*:  
 $\text{Rep-matrix } (\text{take-columns } A \ c) \ j \ i =$   
 $(\text{if } i < c \text{ then } (\text{Rep-matrix } A \ j \ i) \text{ else } 0)$   
 <proof>

**lemma** *Rep-take-rows[simp]*:  
 $\text{Rep-matrix } (\text{take-rows } A \ r) \ j \ i =$   
 $(\text{if } j < r \text{ then } (\text{Rep-matrix } A \ j \ i) \text{ else } 0)$   
 <proof>

**lemma** *Rep-column-of-matrix[simp]*:  
 $\text{Rep-matrix } (\text{column-of-matrix } A \ c) \ j \ i = (\text{if } i = 0 \text{ then } (\text{Rep-matrix } A \ j \ c) \text{ else } 0)$   
 <proof>

**lemma** *Rep-row-of-matrix[simp]*:  
 $\text{Rep-matrix } (\text{row-of-matrix } A \ r) \ j \ i = (\text{if } j = 0 \text{ then } (\text{Rep-matrix } A \ r \ i) \text{ else } 0)$   
 <proof>

**lemma** *column-of-matrix*:  $\text{ncols } A \leq n \implies \text{column-of-matrix } A \ n = 0$   
 <proof>

**lemma** *row-of-matrix*:  $\text{nrows } A \leq n \implies \text{row-of-matrix } A \ n = 0$   
 <proof>

**lemma** *mult-matrix-singleton-right[simp]*:  
**assumes** *prems*:  
 $! x. \text{fmul } x \ 0 = 0$   
 $! x. \text{fmul } 0 \ x = 0$   
 $! x. \text{fadd } 0 \ x = x$   
 $! x. \text{fadd } x \ 0 = x$   
**shows**  $(\text{mult-matrix } \text{fmul } \text{fadd } A \ (\text{singleton-matrix } j \ i \ e)) = \text{apply-matrix } (\% x. \text{fmul } x \ e) \ (\text{move-matrix } (\text{column-of-matrix } A \ j) \ 0 \ (\text{int } i))$   
 <proof>

**lemma** *mult-matrix-ext*:

```

assumes
  eprem:
    ? e. (! a b. a ≠ b ⟶ fmul a e ≠ fmul b e)
and fprems:
  ! a. fmul 0 a = 0
  ! a. fmul a 0 = 0
  ! a. fadd a 0 = a
  ! a. fadd 0 a = a
and contraprems:
  mult-matrix fmul fadd A = mult-matrix fmul fadd B
shows
  A = B
⟨proof⟩

constdefs
  foldmatrix :: ('a ⇒ 'a ⇒ 'a) ⇒ ('a ⇒ 'a ⇒ 'a) ⇒ ('a infmatrix) ⇒ nat ⇒ nat
    ⇒ 'a
  foldmatrix f g A m n == foldseq-transposed g (% j. foldseq f (A j) n) m
  foldmatrix-transposed :: ('a ⇒ 'a ⇒ 'a) ⇒ ('a ⇒ 'a ⇒ 'a) ⇒ ('a infmatrix) ⇒
    nat ⇒ nat ⇒ 'a
  foldmatrix-transposed f g A m n == foldseq g (% j. foldseq-transposed f (A j) n)
    m

lemma foldmatrix-transpose:
assumes
  ! a b c d. g(f a b) (f c d) = f (g a c) (g b d)
shows
  foldmatrix f g A m n = foldmatrix-transposed g f (transpose-infmatrix A) n m
(is ?concl)
⟨proof⟩

lemma foldseq-foldseq:
assumes
  associative f
  associative g
  ! a b c d. g(f a b) (f c d) = f (g a c) (g b d)
shows
  foldseq g (% j. foldseq f (A j) n) m = foldseq f (% j. foldseq g ((transpose-infmatrix
    A) j) m) n
  ⟨proof⟩

lemma mult-n-nrows:
assumes
  ! a. fmul 0 a = 0
  ! a. fmul a 0 = 0
  fadd 0 0 = 0
shows nrows (mult-matrix-n n fmul fadd A B) ≤ nrows A
  ⟨proof⟩

```

**lemma** *mult-n-ncols*:

**assumes**

$! a. \text{fmul } 0 \ a = 0$

$! a. \text{fmul } a \ 0 = 0$

$\text{fadd } 0 \ 0 = 0$

**shows**  $\text{ncols } (\text{mult-matrix-n } n \ \text{fmul } \text{fadd } A \ B) \leq \text{ncols } B$

$\langle \text{proof} \rangle$

**lemma** *mult-nrows*:

**assumes**

$! a. \text{fmul } 0 \ a = 0$

$! a. \text{fmul } a \ 0 = 0$

$\text{fadd } 0 \ 0 = 0$

**shows**  $\text{nrows } (\text{mult-matrix } \text{fmul } \text{fadd } A \ B) \leq \text{nrows } A$

$\langle \text{proof} \rangle$

**lemma** *mult-ncols*:

**assumes**

$! a. \text{fmul } 0 \ a = 0$

$! a. \text{fmul } a \ 0 = 0$

$\text{fadd } 0 \ 0 = 0$

**shows**  $\text{ncols } (\text{mult-matrix } \text{fmul } \text{fadd } A \ B) \leq \text{ncols } B$

$\langle \text{proof} \rangle$

**lemma** *nrows-move-matrix-le*:  $\text{nrows } (\text{move-matrix } A \ j \ i) \leq \text{nat}((\text{int } (\text{nrows } A)) + j)$

$\langle \text{proof} \rangle$

**lemma** *ncols-move-matrix-le*:  $\text{ncols } (\text{move-matrix } A \ j \ i) \leq \text{nat}((\text{int } (\text{ncols } A)) + i)$

$\langle \text{proof} \rangle$

**lemma** *mult-matrix-assoc*:

**assumes** *prems*:

$! a. \text{fmul1 } 0 \ a = 0$

$! a. \text{fmul1 } a \ 0 = 0$

$! a. \text{fmul2 } 0 \ a = 0$

$! a. \text{fmul2 } a \ 0 = 0$

$\text{fadd1 } 0 \ 0 = 0$

$\text{fadd2 } 0 \ 0 = 0$

$! a \ b \ c \ d. \text{fadd2 } (\text{fadd1 } a \ b) (\text{fadd1 } c \ d) = \text{fadd1 } (\text{fadd2 } a \ c) (\text{fadd2 } b \ d)$

*associative fadd1*

*associative fadd2*

$! a \ b \ c. \text{fmul2 } (\text{fmul1 } a \ b) \ c = \text{fmul1 } a \ (\text{fmul2 } b \ c)$

$! a \ b \ c. \text{fmul2 } (\text{fadd1 } a \ b) \ c = \text{fadd1 } (\text{fmul2 } a \ c) (\text{fmul2 } b \ c)$

$! a \ b \ c. \text{fmul1 } c \ (\text{fadd2 } a \ b) = \text{fadd2 } (\text{fmul1 } c \ a) (\text{fmul1 } c \ b)$

**shows**  $\text{mult-matrix } \text{fmul2 } \text{fadd2 } (\text{mult-matrix } \text{fmul1 } \text{fadd1 } A \ B) \ C = \text{mult-matrix } \text{fmul1 } \text{fadd1 } A \ (\text{mult-matrix } \text{fmul2 } \text{fadd2 } B \ C)$  (**is** *?concl*)

$\langle \text{proof} \rangle$

**lemma**

**assumes** *prems*:

! *a*. *fmul1* 0 *a* = 0

! *a*. *fmul1* *a* 0 = 0

! *a*. *fmul2* 0 *a* = 0

! *a*. *fmul2* *a* 0 = 0

*fadd1* 0 0 = 0

*fadd2* 0 0 = 0

! *a b c d*. *fadd2* (*fadd1* *a b*) (*fadd1* *c d*) = *fadd1* (*fadd2* *a c*) (*fadd2* *b d*)

*associative fadd1*

*associative fadd2*

! *a b c*. *fmul2* (*fmul1* *a b*) *c* = *fmul1* *a* (*fmul2* *b c*)

! *a b c*. *fmul2* (*fadd1* *a b*) *c* = *fadd1* (*fmul2* *a c*) (*fmul2* *b c*)

! *a b c*. *fmul1* *c* (*fadd2* *a b*) = *fadd2* (*fmul1* *c a*) (*fmul1* *c b*)

**shows**

(*mult-matrix fmul1 fadd1 A*) o (*mult-matrix fmul2 fadd2 B*) = *mult-matrix fmul2 fadd2 (mult-matrix fmul1 fadd1 A B)*  
 <proof>

**lemma** *mult-matrix-assoc-simple*:

**assumes** *prems*:

! *a*. *fmul* 0 *a* = 0

! *a*. *fmul* *a* 0 = 0

*fadd* 0 0 = 0

*associative fadd*

*commutative fadd*

*associative fmul*

*distributive fmul fadd*

**shows** *mult-matrix fmul fadd (mult-matrix fmul fadd A B) C* = *mult-matrix fmul fadd A (mult-matrix fmul fadd B C)* (is ?concl)  
 <proof>

**lemma** *transpose-apply-matrix*: *f* 0 = 0  $\implies$  *transpose-matrix (apply-matrix f A)*  
 = *apply-matrix f (transpose-matrix A)*  
 <proof>

**lemma** *transpose-combine-matrix*: *f* 0 0 = 0  $\implies$  *transpose-matrix (combine-matrix f A B)*  
 = *combine-matrix f (transpose-matrix A) (transpose-matrix B)*  
 <proof>

**lemma** *Rep-mult-matrix*:

**assumes**

! *a*. *fmul* 0 *a* = 0

! *a*. *fmul* *a* 0 = 0

*fadd* 0 0 = 0

**shows**

*Rep-matrix(mult-matrix fmul fadd A B) j i* =  
*foldseq fadd (% k. fmul (Rep-matrix A j k) (Rep-matrix B k i)) (max (ncols A)*

(*nrows B*))  
 <proof>

**lemma** *transpose-mult-matrix*:

**assumes**

! *a*. *fmul* 0 *a* = 0

! *a*. *fmul* *a* 0 = 0

*fadd* 0 0 = 0

! *x y*. *fmul* *y x* = *fmul* *x y*

**shows**

*transpose-matrix* (*mult-matrix fmul fadd A B*) = *mult-matrix fmul fadd* (*transpose-matrix B*) (*transpose-matrix A*)

<proof>

**lemma** *column-transpose-matrix*: *column-of-matrix* (*transpose-matrix A*) *n* = *transpose-matrix* (*row-of-matrix A n*)

<proof>

**lemma** *take-columns-transpose-matrix*: *take-columns* (*transpose-matrix A*) *n* = *transpose-matrix* (*take-rows A n*)

<proof>

**instantiation** *matrix* :: ({*ord*, *zero*}) *ord*  
**begin**

**definition**

*le-matrix-def*:  $A \leq B \longleftrightarrow (\forall j\ i. \text{Rep-matrix } A\ j\ i \leq \text{Rep-matrix } B\ j\ i)$

**definition**

*less-def*:  $A < (B :: 'a\ matrix) \longleftrightarrow A \leq B \wedge A \neq B$

**instance** <proof>

**end**

**instance** *matrix* :: ({*order*, *zero*}) *order*  
 <proof>

**lemma** *le-apply-matrix*:

**assumes**

*f* 0 = 0

! *x y*. *x* <= *y*  $\longrightarrow$  *f x* <= *f y*

(*a*::('a::({*ord*, *zero*}) *matrix*) <= *b*

**shows**

*apply-matrix f a* <= *apply-matrix f b*

<proof>

**lemma** *le-combine-matrix*:

**assumes**

$f\ 0\ 0 = 0$   
 $! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow f\ a\ c \leq f\ b\ d$   
 $A \leq B$   
 $C \leq D$   
**shows**  
 $combine\_matrix\ f\ A\ C \leq combine\_matrix\ f\ B\ D$   
 $\langle proof \rangle$

**lemma** *le-left-combine-matrix*:  
**assumes**  
 $f\ 0\ 0 = 0$   
 $! a\ b\ c. a \leq b \longrightarrow f\ c\ a \leq f\ c\ b$   
 $A \leq B$   
**shows**  
 $combine\_matrix\ f\ C\ A \leq combine\_matrix\ f\ C\ B$   
 $\langle proof \rangle$

**lemma** *le-right-combine-matrix*:  
**assumes**  
 $f\ 0\ 0 = 0$   
 $! a\ b\ c. a \leq b \longrightarrow f\ a\ c \leq f\ b\ c$   
 $A \leq B$   
**shows**  
 $combine\_matrix\ f\ A\ C \leq combine\_matrix\ f\ B\ C$   
 $\langle proof \rangle$

**lemma** *le-transpose-matrix*:  $(A \leq B) = (transpose\_matrix\ A \leq transpose\_matrix\ B)$   
 $\langle proof \rangle$

**lemma** *le-foldseq*:  
**assumes**  
 $! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow f\ a\ c \leq f\ b\ d$   
 $! i. i \leq n \longrightarrow s\ i \leq t\ i$   
**shows**  
 $foldseq\ f\ s\ n \leq foldseq\ f\ t\ n$   
 $\langle proof \rangle$

**lemma** *le-left-mult*:  
**assumes**  
 $! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow fadd\ a\ c \leq fadd\ b\ d$   
 $! c\ a\ b. 0 \leq c \ \& \ a \leq b \longrightarrow fmul\ c\ a \leq fmul\ c\ b$   
 $! a. fmul\ 0\ a = 0$   
 $! a. fmul\ a\ 0 = 0$   
 $fadd\ 0\ 0 = 0$   
 $0 \leq C$   
 $A \leq B$   
**shows**  
 $mult\_matrix\ fmul\ fadd\ C\ A \leq mult\_matrix\ fmul\ fadd\ C\ B$



$\langle proof \rangle$

**lemma** *le-right-mult*:

**assumes**

$! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow fadd\ a\ c \leq fadd\ b\ d$

$! c\ a\ b. 0 \leq c \ \& \ a \leq b \longrightarrow fmul\ a\ c \leq fmul\ b\ c$

$! a. fmul\ 0\ a = 0$

$! a. fmul\ a\ 0 = 0$

$fadd\ 0\ 0 = 0$

$0 \leq C$

$A \leq B$

**shows**

$mult\text{-}matrix\ fmul\ fadd\ A\ C \leq mult\text{-}matrix\ fmul\ fadd\ B\ C$

$\langle proof \rangle$

**lemma** *spec2*:  $! j\ i. P\ j\ i \Longrightarrow P\ j\ i\ \langle proof \rangle$

**lemma** *neg-imp*:  $(\neg Q \longrightarrow \neg P) \Longrightarrow P \longrightarrow Q\ \langle proof \rangle$

**lemma** *singleton-matrix-le[simp]*:  $(singleton\text{-}matrix\ j\ i\ a \leq singleton\text{-}matrix\ j\ i\ b) = (a \leq (b::'a::\{order\}))$

$\langle proof \rangle$

**lemma** *singleton-le-zero[simp]*:  $(singleton\text{-}matrix\ j\ i\ x \leq 0) = (x \leq (0::'a::\{order,zero\}))$

$\langle proof \rangle$

**lemma** *singleton-ge-zero[simp]*:  $(0 \leq singleton\text{-}matrix\ j\ i\ x) = ((0::'a::\{order,zero\}) \leq x)$

$\langle proof \rangle$

**lemma** *move-matrix-le-zero[simp]*:  $0 \leq j \Longrightarrow 0 \leq i \Longrightarrow (move\text{-}matrix\ A\ j\ i \leq 0) = (A \leq (0::'a::\{order,zero\})\ matrix)$

$\langle proof \rangle$

**lemma** *move-matrix-zero-le[simp]*:  $0 \leq j \Longrightarrow 0 \leq i \Longrightarrow (0 \leq move\text{-}matrix\ A\ j\ i) = ((0::'a::\{order,zero\})\ matrix \leq A)$

$\langle proof \rangle$

**lemma** *move-matrix-le-move-matrix-iff[simp]*:  $0 \leq j \Longrightarrow 0 \leq i \Longrightarrow (move\text{-}matrix\ A\ j\ i \leq move\text{-}matrix\ B\ j\ i) = (A \leq (B::'a::\{order,zero\})\ matrix)$

$\langle proof \rangle$

**end**

**theory** *Matrix*

**imports** *MatrixGeneral*

**begin**

**instantiation** *matrix* :: (*{zero, lattice}*) *lattice*  
**begin**

**definition**  
*inf* = *combine-matrix inf*

**definition**  
*sup* = *combine-matrix sup*

**instance**  
*<proof>*

**end**

**instantiation** *matrix* :: (*{plus, zero}*) *plus*  
**begin**

**definition**  
*plus-matrix-def*:  $A + B = \text{combine-matrix } (op +) A B$

**instance** *<proof>*

**end**

**instantiation** *matrix* :: (*{uminus, zero}*) *uminus*  
**begin**

**definition**  
*minus-matrix-def*:  $- A = \text{apply-matrix } uminus A$

**instance** *<proof>*

**end**

**instantiation** *matrix* :: (*{minus, zero}*) *minus*  
**begin**

**definition**  
*diff-matrix-def*:  $A - B = \text{combine-matrix } (op -) A B$

**instance** *<proof>*

**end**

**instantiation** *matrix* :: (*{plus, times, zero}*) *times*  
**begin**

**definition**  
*times-matrix-def*:  $A * B = \text{mult-matrix } (op *) (op +) A B$

**instance**  $\langle proof \rangle$

**end**

**instantiation**  $matrix :: (ordered-ab-group-add) \text{ abs}$   
**begin**

**definition**

$abs\text{-}matrix\text{-}def: abs (A :: 'a \text{ matrix}) = sup A (- A)$

**instance**  $\langle proof \rangle$

**end**

**instance**  $matrix :: (ordered-ab-group-add) \text{ ordered-ab-group-add-meet}$   
 $\langle proof \rangle$

**instance**  $matrix :: (ordered-ring) \text{ ordered-ring}$   
 $\langle proof \rangle$

**lemma**  $Rep\text{-}matrix\text{-}add[simp]:$

$Rep\text{-}matrix ((a::('a::ordered-ab-group-add) \text{ matrix}) + b) j i = (Rep\text{-}matrix a j i)$   
 $+ (Rep\text{-}matrix b j i)$   
 $\langle proof \rangle$

**lemma**  $Rep\text{-}matrix\text{-}mult: Rep\text{-}matrix ((a::('a::ordered-ring) \text{ matrix}) * b) j i =$   
 $foldseq (op +) (\% k. (Rep\text{-}matrix a j k) * (Rep\text{-}matrix b k i)) (max (ncols a)$   
 $(nrows b))$   
 $\langle proof \rangle$

**lemma**  $apply\text{-}matrix\text{-}add: ! x y. f (x+y) = (f x) + (f y) \implies f 0 = (0::'a) \implies$   
 $apply\text{-}matrix f ((a::('a::ordered-ab-group-add) \text{ matrix}) + b) = (apply\text{-}matrix f a)$   
 $+ (apply\text{-}matrix f b)$   
 $\langle proof \rangle$

**lemma**  $singleton\text{-}matrix\text{-}add: singleton\text{-}matrix j i ((a::('a::ordered-ab-group-add) + b)$   
 $= (singleton\text{-}matrix j i a) + (singleton\text{-}matrix j i b)$   
 $\langle proof \rangle$

**lemma**  $nrows\text{-}mult: nrows ((A::('a::ordered-ring) \text{ matrix}) * B) \leq nrows A$   
 $\langle proof \rangle$

**lemma**  $ncols\text{-}mult: ncols ((A::('a::ordered-ring) \text{ matrix}) * B) \leq ncols B$   
 $\langle proof \rangle$

**definition**

$one\text{-}matrix :: nat \Rightarrow ('a::\{zero,one\}) \text{ matrix}$  **where**  
 $one\text{-}matrix n = Abs\text{-}matrix (\% j i. \text{if } j = i \ \& \ j < n \text{ then } 1 \text{ else } 0)$

**lemma** *Rep-one-matrix[simp]*: *Rep-matrix (one-matrix n) j i = (if (j = i & j < n) then 1 else 0)*  
 <proof>

**lemma** *nrows-one-matrix[simp]*: *nrows ((one-matrix n) :: ('a::zero-neq-one)matrix) = n (is ?r = -)*  
 <proof>

**lemma** *ncols-one-matrix[simp]*: *ncols ((one-matrix n) :: ('a::zero-neq-one)matrix) = n (is ?r = -)*  
 <proof>

**lemma** *one-matrix-mult-right[simp]*: *ncols A <= n  $\implies$  (A::('a::{lordered-ring, ring-1}) matrix) \* (one-matrix n) = A*  
 <proof>

**lemma** *one-matrix-mult-left[simp]*: *nrows A <= n  $\implies$  (one-matrix n) \* A = (A::('a::{lordered-ring, ring-1}) matrix)*  
 <proof>

**lemma** *transpose-matrix-mult*: *transpose-matrix ((A::('a::{lordered-ring, comm-ring}) matrix)\*B) = (transpose-matrix B) \* (transpose-matrix A)*  
 <proof>

**lemma** *transpose-matrix-add*: *transpose-matrix ((A::('a::lordered-ab-group-add) matrix)+B) = transpose-matrix A + transpose-matrix B*  
 <proof>

**lemma** *transpose-matrix-diff*: *transpose-matrix ((A::('a::lordered-ab-group-add) matrix)-B) = transpose-matrix A - transpose-matrix B*  
 <proof>

**lemma** *transpose-matrix-minus*: *transpose-matrix (-(A::('a::lordered-ring) matrix)) = - transpose-matrix (A::('a::lordered-ring) matrix)*  
 <proof>

**constdefs**

*right-inverse-matrix* :: ('a::{lordered-ring, ring-1}) matrix  $\Rightarrow$  'a matrix  $\Rightarrow$  bool  
*right-inverse-matrix* A X == (A \* X = one-matrix (max (nrows A) (ncols X)))  
 $\wedge$  nrows X  $\leq$  ncols A  
*left-inverse-matrix* :: ('a::{lordered-ring, ring-1}) matrix  $\Rightarrow$  'a matrix  $\Rightarrow$  bool  
*left-inverse-matrix* A X == (X \* A = one-matrix (max(nrows X) (ncols A)))  $\wedge$   
 ncols X  $\leq$  nrows A  
*inverse-matrix* :: ('a::{lordered-ring, ring-1}) matrix  $\Rightarrow$  'a matrix  $\Rightarrow$  bool  
*inverse-matrix* A X == (right-inverse-matrix A X)  $\wedge$  (left-inverse-matrix A X)

**lemma** *right-inverse-matrix-dim*: *right-inverse-matrix A X  $\implies$  nrows A = ncols X*

$\langle \text{proof} \rangle$

**lemma** *left-inverse-matrix-dim*: *left-inverse-matrix*  $A$   $Y \implies \text{ncols } A = \text{nrows } Y$   
 $\langle \text{proof} \rangle$

**lemma** *left-right-inverse-matrix-unique*:  
  **assumes** *left-inverse-matrix*  $A$   $Y$  *right-inverse-matrix*  $A$   $X$   
  **shows**  $X = Y$   
 $\langle \text{proof} \rangle$

**lemma** *inverse-matrix-inject*:  $\llbracket \text{inverse-matrix } A \ X; \text{inverse-matrix } A \ Y \rrbracket \implies X = Y$   
 $\langle \text{proof} \rangle$

**lemma** *one-matrix-inverse*: *inverse-matrix* (*one-matrix*  $n$ ) (*one-matrix*  $n$ )  
 $\langle \text{proof} \rangle$

**lemma** *zero-imp-mult-zero*:  $(a::'a::\text{ring}) = 0 \mid b = 0 \implies a * b = 0$   
 $\langle \text{proof} \rangle$

**lemma** *Rep-matrix-zero-imp-mult-zero*:  
   $! j \ i \ k. (\text{Rep-matrix } A \ j \ k = 0) \mid (\text{Rep-matrix } B \ k \ i) = 0 \implies A * B =$   
   $(0::('a::\text{lordered-ring}) \text{ matrix})$   
 $\langle \text{proof} \rangle$

**lemma** *add-nrows*:  $\text{nrows } (A::('a::\text{comm-monoid-add}) \text{ matrix}) \leq u \implies \text{nrows } B \leq u \implies \text{nrows } (A + B) \leq u$   
 $\langle \text{proof} \rangle$

**lemma** *move-matrix-row-mult*: *move-matrix*  $((A::('a::\text{lordered-ring}) \text{ matrix}) * B)$   
 $j \ 0 = (\text{move-matrix } A \ j \ 0) * B$   
 $\langle \text{proof} \rangle$

**lemma** *move-matrix-col-mult*: *move-matrix*  $((A::('a::\text{lordered-ring}) \text{ matrix}) * B)$   
 $0 \ i = A * (\text{move-matrix } B \ 0 \ i)$   
 $\langle \text{proof} \rangle$

**lemma** *move-matrix-add*:  $((\text{move-matrix } (A + B) \ j \ i)::('a::\text{lordered-ab-group-add}) \text{ matrix})) = (\text{move-matrix } A \ j \ i) + (\text{move-matrix } B \ j \ i)$   
 $\langle \text{proof} \rangle$

**lemma** *move-matrix-mult*: *move-matrix*  $((A::('a::\text{lordered-ring}) \text{ matrix}) * B) \ j \ i = (\text{move-matrix } A \ j \ 0) * (\text{move-matrix } B \ 0 \ i)$   
 $\langle \text{proof} \rangle$

**constdefs**  
  *scalar-mult*  $:: ('a::\text{lordered-ring}) \Rightarrow 'a \text{ matrix} \Rightarrow 'a \text{ matrix}$   
  *scalar-mult*  $a \ m == \text{apply-matrix } (\text{op } * \ a) \ m$

**lemma** *scalar-mult-zero[simp]*:  $\text{scalar-mult } y \ 0 = 0$   
 $\langle \text{proof} \rangle$

**lemma** *scalar-mult-add*:  $\text{scalar-mult } y \ (a+b) = (\text{scalar-mult } y \ a) + (\text{scalar-mult } y \ b)$   
 $\langle \text{proof} \rangle$

**lemma** *Rep-scalar-mult[simp]*:  $\text{Rep-matrix } (\text{scalar-mult } y \ a) \ j \ i = y * (\text{Rep-matrix } a \ j \ i)$   
 $\langle \text{proof} \rangle$

**lemma** *scalar-mult-singleton[simp]*:  $\text{scalar-mult } y \ (\text{singleton-matrix } j \ i \ x) = \text{singleton-matrix } j \ i \ (y * x)$   
 $\langle \text{proof} \rangle$

**lemma** *Rep-minus[simp]*:  $\text{Rep-matrix } (-(A:::\text{lordered-ab-group-add})) \ x \ y = -(\text{Rep-matrix } A \ x \ y)$   
 $\langle \text{proof} \rangle$

**lemma** *Rep-abs[simp]*:  $\text{Rep-matrix } (\text{abs } (A:::\text{lordered-ring})) \ x \ y = \text{abs } (\text{Rep-matrix } A \ x \ y)$   
 $\langle \text{proof} \rangle$

**end**

**theory** *LP*  
**imports** *Main*  
**begin**

**lemma** *linprog-dual-estimate*:  
**assumes**  
 $A * x \leq (b::'\text{a}::\text{lordered-ring})$   
 $0 \leq y$   
 $\text{abs } (A - A') \leq \delta A$   
 $b \leq b'$   
 $\text{abs } (c - c') \leq \delta c$   
 $\text{abs } x \leq r$   
**shows**  
 $c * x \leq y * b' + (y * \delta A + \text{abs } (y * A' - c') + \delta c) * r$   
 $\langle \text{proof} \rangle$

**lemma** *le-ge-imp-abs-diff-1*:  
**assumes**  
 $A1 \leq (A::'\text{a}::\text{lordered-ring})$   
 $A \leq A2$   
**shows**  $\text{abs } (A - A1) \leq A2 - A1$   
 $\langle \text{proof} \rangle$

**lemma** *mult-le-prts*:

**assumes**

$a1 \leq (a::'a::\text{ordered-ring})$

$a \leq a2$

$b1 \leq b$

$b \leq b2$

**shows**

$a * b \leq \text{pprt } a2 * \text{pprt } b2 + \text{pprt } a1 * \text{nprt } b2 + \text{nprt } a2 * \text{pprt } b1 + \text{nprt } a1$   
 $* \text{nprt } b1$   
 $\langle \text{proof} \rangle$

**lemma** *mult-le-dual-prts*:

**assumes**

$A * x \leq (b::'a::\text{ordered-ring})$

$0 \leq y$

$A1 \leq A$

$A \leq A2$

$c1 \leq c$

$c \leq c2$

$r1 \leq x$

$x \leq r2$

**shows**

$c * x \leq y * b + (\text{let } s1 = c1 - y * A2; s2 = c2 - y * A1 \text{ in } \text{pprt } s2 * \text{pprt } r2$   
 $+ \text{pprt } s1 * \text{nprt } r2 + \text{nprt } s2 * \text{pprt } r1 + \text{nprt } s1 * \text{nprt } r1)$   
 $(\text{is } - \leq - + ?C)$   
 $\langle \text{proof} \rangle$

**end**

**theory** *SparseMatrix* **imports** *Matrix LP* **begin**

**types**

$'a \text{ svec} = (\text{nat} * 'a) \text{ list}$

$'a \text{ spmat} = ('a \text{ svec}) \text{ svec}$

**consts**

$\text{sparse-row-vector} :: ('a::\text{ordered-ring}) \text{ svec} \Rightarrow 'a \text{ matrix}$

$\text{sparse-row-matrix} :: ('a::\text{ordered-ring}) \text{ spmat} \Rightarrow 'a \text{ matrix}$

**defs**

$\text{sparse-row-vector-def} : \text{sparse-row-vector } arr == \text{foldl } (\% m \ x. m + (\text{singleton-matrix } 0 \ (\text{fst } x) \ (\text{snd } x))) \ 0 \ arr$

$\text{sparse-row-matrix-def} : \text{sparse-row-matrix } arr == \text{foldl } (\% m \ r. m + (\text{move-matrix } (\text{sparse-row-vector } (\text{snd } r)) \ (\text{int } (\text{fst } r)) \ 0)) \ 0 \ arr$

**lemma** *sparse-row-vector-empty[simp]*:  $\text{sparse-row-vector } [] = 0$   
 $\langle \text{proof} \rangle$

**lemma** *sparse-row-matrix-empty*[simp]: *sparse-row-matrix* [] = 0  
 ⟨proof⟩

**lemma** *foldl-distrstart*[rule-format]: ! a x y. (f (g x y) a = g x (f y a)) ==> ! x y.  
 (foldl f (g x y) l = g x (foldl f y l))  
 ⟨proof⟩

**lemma** *sparse-row-vector-cons*[simp]: *sparse-row-vector* (a#arr) = (*singleton-matrix* 0 (fst a) (snd a)) + (*sparse-row-vector* arr)  
 ⟨proof⟩

**lemma** *sparse-row-vector-append*[simp]: *sparse-row-vector* (a @ b) = (*sparse-row-vector* a) + (*sparse-row-vector* b)  
 ⟨proof⟩

**lemma** *nrows-spvec*[simp]: *nrows* (*sparse-row-vector* x) <= (Suc 0)  
 ⟨proof⟩

**lemma** *sparse-row-matrix-cons*: *sparse-row-matrix* (a#arr) = ((*move-matrix* (*sparse-row-vector* (snd a)) (int (fst a)) 0)) + *sparse-row-matrix* arr  
 ⟨proof⟩

**lemma** *sparse-row-matrix-append*: *sparse-row-matrix* (arr@brr) = (*sparse-row-matrix* arr) + (*sparse-row-matrix* brr)  
 ⟨proof⟩

**consts**  
*sorted-spvec* :: 'a spvec => bool  
*sorted-spmat* :: 'a spat => bool

**primrec**  
*sorted-spmat* [] = True  
*sorted-spmat* (a#as) = ((*sorted-spvec* (snd a)) & (*sorted-spmat* as))

**primrec**  
*sorted-spvec* [] = True  
*sorted-spvec-step*: *sorted-spvec* (a#as) = (case as of [] => True | b#bs => ((fst a < fst b) & (*sorted-spvec* as)))

**declare** *sorted-spvec.simps* [simp del]

**lemma** *sorted-spvec-empty*[simp]: *sorted-spvec* [] = True  
 ⟨proof⟩

**lemma** *sorted-spvec-cons1*: *sorted-spvec* (a#as) ==> *sorted-spvec* as  
 ⟨proof⟩

**lemma** *sorted-spvec-cons2*: *sorted-spvec* (a#b#t) ==> *sorted-spvec* (a#t)



$\langle \text{proof} \rangle$

**lemma** *sorted-spvec-cons3*:  $\text{sorted-spvec}(a \# b \# t) \implies \text{fst } a < \text{fst } b$   
 $\langle \text{proof} \rangle$

**lemma** *sorted-sparse-row-vector-zero*[*rule-format*]:  $m \leq n \longrightarrow \text{sorted-spvec } ((n, a) \# \text{arr})$   
 $\longrightarrow \text{Rep-matrix } (\text{sparse-row-vector } \text{arr}) \ j \ m = 0$   
 $\langle \text{proof} \rangle$

**lemma** *sorted-sparse-row-matrix-zero*[*rule-format*]:  $m \leq n \longrightarrow \text{sorted-spvec } ((n, a) \# \text{arr})$   
 $\longrightarrow \text{Rep-matrix } (\text{sparse-row-matrix } \text{arr}) \ m \ j = 0$   
 $\langle \text{proof} \rangle$

**consts**

*abs-spvec* ::  $('a :: \text{lordered-ring}) \text{ spvec} \Rightarrow 'a \text{ spvec}$   
*minus-spvec* ::  $('a :: \text{lordered-ring}) \text{ spvec} \Rightarrow 'a \text{ spvec}$   
*smult-spvec* ::  $('a :: \text{lordered-ring}) \Rightarrow 'a \text{ spvec} \Rightarrow 'a \text{ spvec}$   
*addmult-spvec* ::  $('a :: \text{lordered-ring}) * 'a \text{ spvec} * 'a \text{ spvec} \Rightarrow 'a \text{ spvec}$

**primrec**

*minus-spvec* [] = []  
*minus-spvec* (a # as) = (fst a, -(snd a)) # (minus-spvec as)

**primrec**

*abs-spvec* [] = []  
*abs-spvec* (a # as) = (fst a, abs (snd a)) # (abs-spvec as)

**lemma** *sparse-row-vector-minus*:

*sparse-row-vector* (minus-spvec v) = - (sparse-row-vector v)  
 $\langle \text{proof} \rangle$

**lemma** *sparse-row-vector-abs*:

*sorted-spvec* v  $\implies$  *sparse-row-vector* (abs-spvec v) = abs (sparse-row-vector v)  
 $\langle \text{proof} \rangle$

**lemma** *sorted-spvec-minus-spvec*:

*sorted-spvec* v  $\implies$  *sorted-spvec* (minus-spvec v)  
 $\langle \text{proof} \rangle$

**lemma** *sorted-spvec-abs-spvec*:

*sorted-spvec* v  $\implies$  *sorted-spvec* (abs-spvec v)  
 $\langle \text{proof} \rangle$

**defs**

*smult-spvec-def*: *smult-spvec* y arr == map (% a. (fst a, y \* snd a)) arr

**lemma** *smult-spvec-empty*[*simp*]: *smult-spvec* y [] = []  
 $\langle \text{proof} \rangle$

**lemma** *smult-spvec-cons*:  $smult\text{-}spvec\ y\ (a\#arr) = (fst\ a,\ y * (snd\ a)) \# (smult\text{-}spvec\ y\ arr)$   
 <proof>

**recdef** *addmult-spvec measure* (% (y, a, b). length a + (length b))  
 addmult-spvec (y, arr, []) = arr  
 addmult-spvec (y, [], brr) = smult-spvec y brr  
 addmult-spvec (y, a#arr, b#brr) = (  
   if (fst a) < (fst b) then (a#(addmult-spvec (y, arr, b#brr)))  
   else (if (fst b < fst a) then ((fst b, y \* (snd b))#(addmult-spvec (y, a#arr, brr)))  
   else ((fst a, (snd a)+ y\*(snd b))#(addmult-spvec (y, arr, brr))))))

**lemma** *addmult-spvec-empty1[simp]*:  $addmult\text{-}spvec\ (y,\ [],\ a) = smult\text{-}spvec\ y\ a$   
 <proof>

**lemma** *addmult-spvec-empty2[simp]*:  $addmult\text{-}spvec\ (y,\ a,\ []) = a$   
 <proof>

**lemma** *sparse-row-vector-map*:  $(! x\ y.\ f\ (x+y) = (f\ x) + (f\ y)) \implies (f::'a \Rightarrow ('a::lordered-ring))$   
 $0 = 0 \implies$   
 sparse-row-vector (map (% x. (fst x, f (snd x))) a) = apply-matrix f (sparse-row-vector a)  
 <proof>

**lemma** *sparse-row-vector-smult*:  $sparse\text{-}row\text{-}vector\ (smult\text{-}spvec\ y\ a) = scalar\text{-}mult\ y\ (sparse\text{-}row\text{-}vector\ a)$   
 <proof>

**lemma** *sparse-row-vector-addmult-spvec*:  $sparse\text{-}row\text{-}vector\ (addmult\text{-}spvec\ (y::'a::lordered\text{-}ring,\ a,\ b)) =$   
 $(sparse\text{-}row\text{-}vector\ a) + (scalar\text{-}mult\ y\ (sparse\text{-}row\text{-}vector\ b))$   
 <proof>

**lemma** *sorted-smult-spvec[rule-format]*:  $sorted\text{-}spvec\ a \implies sorted\text{-}spvec\ (smult\text{-}spvec\ y\ a)$   
 <proof>

**lemma** *sorted-spvec-addmult-spvec-helper*:  $\llbracket sorted\text{-}spvec\ (addmult\text{-}spvec\ (y,\ (a,\ b)\ \# arr,\ brr)); aa < a; sorted\text{-}spvec\ ((a,\ b)\ \# arr);$   
 $sorted\text{-}spvec\ ((aa,\ ba)\ \# brr) \rrbracket \implies sorted\text{-}spvec\ ((aa,\ y * ba)\ \# addmult\text{-}spvec\ (y,\ (a,\ b)\ \# arr,\ brr))$   
 <proof>

**lemma** *sorted-spvec-addmult-spvec-helper2*:  
 $\llbracket sorted\text{-}spvec\ (addmult\text{-}spvec\ (y,\ arr,\ (aa,\ ba)\ \# brr)); a < aa; sorted\text{-}spvec\ ((a,\ b)\ \# arr); sorted\text{-}spvec\ ((aa,\ ba)\ \# brr) \rrbracket$   
 $\implies sorted\text{-}spvec\ ((a,\ b)\ \# addmult\text{-}spvec\ (y,\ arr,\ (aa,\ ba)\ \# brr))$   
 <proof>

**lemma** *sorted-spvec-addmult-spvec-helper3*[rule-format]:

*sorted-spvec* (*addmult-spvec* (*y*, *arr*, *brr*))  $\longrightarrow$  *sorted-spvec* ((*aa*, *b*) # *arr*)  $\longrightarrow$   
*sorted-spvec* ((*aa*, *ba*) # *brr*)  
 $\longrightarrow$  *sorted-spvec* ((*aa*, *b* + *y* \* *ba*) # (*addmult-spvec* (*y*, *arr*, *brr*)))  
 <proof>

**lemma** *sorted-addmult-spvec*[rule-format]: *sorted-spvec* *a*  $\longrightarrow$  *sorted-spvec* *b*  $\longrightarrow$   
*sorted-spvec* (*addmult-spvec* (*y*, *a*, *b*))  
 <proof>

**consts**

*mult-spvec-spmat* :: ('a::lordered-ring) *spvec* \* 'a *spvec* \* 'a *spmat*  $\Rightarrow$  'a *spvec*

**recdef** *mult-spvec-spmat* *measure* (% (*c*, *arr*, *brr*). (length *arr*) + (length *brr*))

*mult-spvec-spmat* (*c*, [], *brr*) = *c*  
*mult-spvec-spmat* (*c*, *arr*, []) = *c*  
*mult-spvec-spmat* (*c*, *a*#*arr*, *b*#*brr*) = (  
 if ((fst *a*) < (fst *b*)) then (*mult-spvec-spmat* (*c*, *arr*, *b*#*brr*))  
 else if ((fst *b*) < (fst *a*)) then (*mult-spvec-spmat* (*c*, *a*#*arr*, *brr*))  
 else (*mult-spvec-spmat* (*addmult-spvec* (snd *a*, *c*, snd *b*), *arr*, *brr*)))

**lemma** *sparse-row-mult-spvec-spmat*[rule-format]: *sorted-spvec* (*a*::('a::lordered-ring)  
*spvec*)  $\longrightarrow$  *sorted-spvec* *B*  $\longrightarrow$

*sparse-row-vector* (*mult-spvec-spmat* (*c*, *a*, *B*)) = (*sparse-row-vector* *c*) + (*sparse-row-vector*  
*a*) \* (*sparse-row-matrix* *B*)  
 <proof>

**lemma** *sorted-mult-spvec-spmat*[rule-format]:

*sorted-spvec* (*c*::('a::lordered-ring) *spvec*)  $\longrightarrow$  *sorted-spmat* *B*  $\longrightarrow$  *sorted-spvec*  
(*mult-spvec-spmat* (*c*, *a*, *B*))  
 <proof>

**consts**

*mult-spmat* :: ('a::lordered-ring) *spmat*  $\Rightarrow$  'a *spmat*  $\Rightarrow$  'a *spmat*

**primrec**

*mult-spmat* [] *A* = []  
*mult-spmat* (*a*#*as*) *A* = (fst *a*, *mult-spvec-spmat* ([], snd *a*, *A*))#(*mult-spmat* *as*  
*A*)

**lemma** *sparse-row-mult-spmat*[rule-format]:

*sorted-spmat* *A*  $\longrightarrow$  *sorted-spvec* *B*  $\longrightarrow$  *sparse-row-matrix* (*mult-spmat* *A* *B*) =  
(*sparse-row-matrix* *A*) \* (*sparse-row-matrix* *B*)  
 <proof>

**lemma** *sorted-spvec-mult-spmat*[rule-format]:

*sorted-spvec* (*A*::('a::lordered-ring) *spmat*)  $\longrightarrow$  *sorted-spvec* (*mult-spmat* *A* *B*)  
 <proof>

**lemma** *sorted-spmat-mult-spmat*[rule-format]:  
 $\text{sorted-spmat } (B::('a::\text{ordered-ring}) \text{ spat}) \longrightarrow \text{sorted-spmat } (\text{mult-spmat } A \ B)$   
 <proof>

**consts**

$\text{add-spvec} :: ('a::\text{ordered-ab-group-add}) \text{ spvec} * 'a \text{ spvec} \Rightarrow 'a \text{ spvec}$   
 $\text{add-spmat} :: ('a::\text{ordered-ab-group-add}) \text{ spat} * 'a \text{ spat} \Rightarrow 'a \text{ spat}$

**recdef** *add-spvec measure* (% (a, b). length a + (length b))  
 $\text{add-spvec } (arr, []) = arr$   
 $\text{add-spvec } ([], brr) = brr$   
 $\text{add-spvec } (a\#arr, b\#brr) = ($   
 if (fst a) < (fst b) then (a#(add-spvec (arr, b#brr)))  
 else (if (fst b < fst a) then (b#(add-spvec (a#arr, brr)))  
 else ((fst a, (snd a)+(snd b))#(add-spvec (arr,brr))))

**lemma** *add-spvec-empty1*[simp]:  $\text{add-spvec } ([], a) = a$   
 <proof>

**lemma** *add-spvec-empty2*[simp]:  $\text{add-spvec } (a, []) = a$   
 <proof>

**lemma** *sparse-row-vector-add*:  $\text{sparse-row-vector } (\text{add-spvec } (a,b)) = (\text{sparse-row-vector } a) + (\text{sparse-row-vector } b)$   
 <proof>

**recdef** *add-spmat measure* (% (A,B). (length A)+(length B))  
 $\text{add-spmat } ([], bs) = bs$   
 $\text{add-spmat } (as, []) = as$   
 $\text{add-spmat } (a\#as, b\#bs) = ($   
 if fst a < fst b then  
 (a#(add-spmat (as, b#bs)))  
 else (if fst b < fst a then  
 (b#(add-spmat (a#as, bs)))  
 else  
 ((fst a, add-spvec (snd a, snd b))#(add-spmat (as, bs))))

**lemma** *sparse-row-add-spmat*:  $\text{sparse-row-matrix } (\text{add-spmat } (A, B)) = (\text{sparse-row-matrix } A) + (\text{sparse-row-matrix } B)$   
 <proof>

**lemma** *sorted-add-spvec-helper1*[rule-format]:  $\text{add-spvec } ((a,b)\#arr, brr) = (ab,$   
 $bb) \# \text{list} \longrightarrow (ab = a \mid (brr \neq [] \ \& \ ab = \text{fst } (\text{hd } brr)))$   
 <proof>

**lemma** *sorted-add-spmat-helper1*[rule-format]:  $\text{add-spmat } ((a,b)\#arr, brr) = (ab,$   
 $bb) \# \text{list} \longrightarrow (ab = a \mid (brr \neq [] \ \& \ ab = \text{fst } (\text{hd } brr)))$   
 <proof>

**lemma** *sorted-add-spvec-helper*[rule-format]:  $\text{add-spvec } (arr, brr) = (ab, bb) \# list \longrightarrow ((arr \neq [] \ \& \ ab = fst \ (hd \ arr)) \mid (brr \neq [] \ \& \ ab = fst \ (hd \ brr)))$   
 <proof>

**lemma** *sorted-add-spmat-helper*[rule-format]:  $\text{add-spmat } (arr, brr) = (ab, bb) \# list \longrightarrow ((arr \neq [] \ \& \ ab = fst \ (hd \ arr)) \mid (brr \neq [] \ \& \ ab = fst \ (hd \ brr)))$   
 <proof>

**lemma** *add-spvec-commute*:  $\text{add-spvec } (a, b) = \text{add-spvec } (b, a)$   
 <proof>

**lemma** *add-spmat-commute*:  $\text{add-spmat } (a, b) = \text{add-spmat } (b, a)$   
 <proof>

**lemma** *sorted-add-spvec-helper2*:  $\text{add-spvec } ((a,b)\#arr, brr) = (ab, bb) \# list \implies aa < a \implies \text{sorted-spvec } ((aa, ba) \# brr) \implies aa < ab$   
 <proof>

**lemma** *sorted-add-spmat-helper2*:  $\text{add-spmat } ((a,b)\#arr, brr) = (ab, bb) \# list \implies aa < a \implies \text{sorted-spmat } ((aa, ba) \# brr) \implies aa < ab$   
 <proof>

**lemma** *sorted-spvec-add-spvec*[rule-format]:  $\text{sorted-spvec } a \longrightarrow \text{sorted-spvec } b \longrightarrow \text{sorted-spvec } (\text{add-spvec } (a, b))$   
 <proof>

**lemma** *sorted-spvec-add-spmat*[rule-format]:  $\text{sorted-spvec } A \longrightarrow \text{sorted-spvec } B \longrightarrow \text{sorted-spmat } (\text{add-spmat } (A, B))$   
 <proof>

**lemma** *sorted-spmat-add-spmat*[rule-format]:  $\text{sorted-spmat } A \longrightarrow \text{sorted-spmat } B \longrightarrow \text{sorted-spmat } (\text{add-spmat } (A, B))$   
 <proof>

**consts**

$\text{le-spvec} :: ('a::\text{ordered-ab-group-add}) \text{ spvec} * 'a \text{ spvec} \Rightarrow \text{bool}$   
 $\text{le-spmat} :: ('a::\text{ordered-ab-group-add}) \text{ spat} * 'a \text{ spat} \Rightarrow \text{bool}$

**recdef** *le-spvec measure* (% (a,b). (length a) + (length b))

$\text{le-spvec } ([], []) = \text{True}$   
 $\text{le-spvec } (a\#as, []) = ((snd \ a \leq 0) \ \& \ (\text{le-spvec } (as, [])))$   
 $\text{le-spvec } ([], b\#bs) = ((0 \leq snd \ b) \ \& \ (\text{le-spvec } ([], bs)))$   
 $\text{le-spmat } (a\#as, b\#bs) =$   
 if (fst a < fst b) then  
 ((snd a <= 0) & (le-spmat (as, b#bs)))  
 else if (fst b < fst a) then  
 ((0 <= snd b) & (le-spmat (a#as, bs)))  
 else

```

((snd a <= snd b) & (le-spvec (as, bs))))

recdef le-spmat measure (% (a,b). (length a) + (length b))
  le-spmat ([], []) = True
  le-spmat (a#as, []) = (le-spvec (snd a, []) & (le-spmat (as, [])))
  le-spmat ([], b#bs) = (le-spvec ([], snd b) & (le-spmat ([], bs)))
  le-spmat (a#as, b#bs) = (
    if fst a < fst b then
      (le-spvec(snd a,[]) & le-spmat(as, b#bs))
    else (if (fst b < fst a) then
      (le-spvec([], snd b) & le-spmat(a#as, bs))
    else
      (le-spvec(snd a, snd b) & le-spmat (as, bs))))

constdefs
  disj-matrices :: ('a::zero) matrix => 'a matrix => bool
  disj-matrices A B == (! j i. (Rep-matrix A j i ≠ 0) → (Rep-matrix B j i = 0)) & (! j i. (Rep-matrix B j i ≠ 0) → (Rep-matrix A j i = 0))

declare [[simp-depth-limit = 6]]

lemma disj-matrices-contr1: disj-matrices A B ==> Rep-matrix A j i ≠ 0 ==>
  Rep-matrix B j i = 0
  <proof>

lemma disj-matrices-contr2: disj-matrices A B ==> Rep-matrix B j i ≠ 0 ==>
  Rep-matrix A j i = 0
  <proof>

lemma disj-matrices-add: disj-matrices A B ==> disj-matrices C D ==> disj-matrices
  A D ==> disj-matrices B C ==>
  (A + B <= C + D) = (A <= C & B <= (D::('a::lordered-ab-group-add)
  matrix))
  <proof>

lemma disj-matrices-zero1[simp]: disj-matrices 0 B
  <proof>

lemma disj-matrices-zero2[simp]: disj-matrices A 0
  <proof>

lemma disj-matrices-commute: disj-matrices A B = disj-matrices B A
  <proof>

lemma disj-matrices-add-le-zero: disj-matrices A B ==>
  (A + B <= 0) = (A <= 0 & (B::('a::lordered-ab-group-add) matrix) <= 0)
  <proof>

```

**lemma** *disj-matrices-add-zero-le*:  $\text{disj-matrices } A \ B \implies$   
 $(0 \leq A + B) = (0 \leq A \ \& \ 0 \leq (B::('a::\text{ordered-ab-group-add}) \text{ matrix}))$   
 $\langle \text{proof} \rangle$

**lemma** *disj-matrices-add-x-le*:  $\text{disj-matrices } A \ B \implies \text{disj-matrices } B \ C \implies$   
 $(A \leq B + C) = (A \leq C \ \& \ 0 \leq (B::('a::\text{ordered-ab-group-add}) \text{ matrix}))$   
 $\langle \text{proof} \rangle$

**lemma** *disj-matrices-add-le-x*:  $\text{disj-matrices } A \ B \implies \text{disj-matrices } B \ C \implies$   
 $(B + A \leq C) = (A \leq C \ \& \ (B::('a::\text{ordered-ab-group-add}) \text{ matrix}) \leq 0)$   
 $\langle \text{proof} \rangle$

**lemma** *disj-sparse-row-singleton*:  $i \leq j \implies \text{sorted-spvec}((j,y)\#v) \implies \text{disj-matrices}$   
 $(\text{sparse-row-vector } v) \ (\text{singleton-matrix } 0 \ i \ x)$   
 $\langle \text{proof} \rangle$

**lemma** *disj-matrices-x-add*:  $\text{disj-matrices } A \ B \implies \text{disj-matrices } A \ C \implies \text{disj-matrices}$   
 $(A::('a::\text{ordered-ab-group-add}) \text{ matrix}) \ (B+C)$   
 $\langle \text{proof} \rangle$

**lemma** *disj-matrices-add-x*:  $\text{disj-matrices } A \ B \implies \text{disj-matrices } A \ C \implies \text{disj-matrices}$   
 $(B+C) \ (A::('a::\text{ordered-ab-group-add}) \text{ matrix})$   
 $\langle \text{proof} \rangle$

**lemma** *disj-singleton-matrices[simp]*:  $\text{disj-matrices} \ (\text{singleton-matrix } j \ i \ x) \ (\text{singleton-matrix}$   
 $u \ v \ y) = (j \neq u \mid i \neq v \mid x = 0 \mid y = 0)$   
 $\langle \text{proof} \rangle$

**lemma** *disj-move-sparse-vec-mat[simplified disj-matrices-commute]*:  
 $j \leq a \implies \text{sorted-spvec}((a,c)\#as) \implies \text{disj-matrices} \ (\text{move-matrix} \ (\text{sparse-row-vector}$   
 $b) \ (\text{int } j) \ i) \ (\text{sparse-row-matrix } as)$   
 $\langle \text{proof} \rangle$

**lemma** *disj-move-sparse-row-vector-twice*:  
 $j \neq u \implies \text{disj-matrices} \ (\text{move-matrix} \ (\text{sparse-row-vector } a) \ j \ i) \ (\text{move-matrix}$   
 $(\text{sparse-row-vector } b) \ u \ v)$   
 $\langle \text{proof} \rangle$

**lemma** *le-spvec-iff-sparse-row-le[rule-format]*:  $(\text{sorted-spvec } a) \longrightarrow (\text{sorted-spvec}$   
 $b) \longrightarrow (\text{le-spvec } (a,b)) = (\text{sparse-row-vector } a \leq \text{sparse-row-vector } b)$   
 $\langle \text{proof} \rangle$

**lemma** *le-spvec-empty2-sparse-row[rule-format]*:  $(\text{sorted-spvec } b) \longrightarrow (\text{le-spvec } (b,[]))$   
 $= (\text{sparse-row-vector } b \leq 0)$   
 $\langle \text{proof} \rangle$

**lemma** *le-spvec-empty1-sparse-row[rule-format]*:  $(\text{sorted-spvec } b) \longrightarrow (\text{le-spvec } ([],b))$   
 $= (0 \leq \text{sparse-row-vector } b)$   
 $\langle \text{proof} \rangle$

**lemma** *le-spmat-iff-sparse-row-le*[*rule-format*]: (*sorted-spvec* *A*)  $\longrightarrow$  (*sorted-spmat* *A*)  $\longrightarrow$  (*sorted-spvec* *B*)  $\longrightarrow$  (*sorted-spmat* *B*)  $\longrightarrow$   
*le-spmat*(*A*, *B*) = (*sparse-row-matrix* *A* <= *sparse-row-matrix* *B*)  
 <proof>

**declare** [[*simp-depth-limit* = 999]]

**consts**

*abs-spmat* :: ('a::lordered-ring) *spmat*  $\Rightarrow$  'a *spmat*  
*minus-spmat* :: ('a::lordered-ring) *spmat*  $\Rightarrow$  'a *spmat*

**primrec**

*abs-spmat* [] = []  
*abs-spmat* (*a*#*as*) = (*fst* *a*, *abs-spvec* (*snd* *a*))#(*abs-spmat* *as*)

**primrec**

*minus-spmat* [] = []  
*minus-spmat* (*a*#*as*) = (*fst* *a*, *minus-spvec* (*snd* *a*))#(*minus-spmat* *as*)

**lemma** *sparse-row-matrix-minus*:

*sparse-row-matrix* (*minus-spmat* *A*) = - (*sparse-row-matrix* *A*)  
 <proof>

**lemma** *Rep-sparse-row-vector-zero*: *x*  $\neq$  0  $\implies$  *Rep-matrix* (*sparse-row-vector* *v*)  
*x* *y* = 0  
 <proof>

**lemma** *sparse-row-matrix-abs*:

*sorted-spvec* *A*  $\implies$  *sorted-spmat* *A*  $\implies$  *sparse-row-matrix* (*abs-spmat* *A*) = *abs*  
 (*sparse-row-matrix* *A*)  
 <proof>

**lemma** *sorted-spvec-minus-spmat*: *sorted-spvec* *A*  $\implies$  *sorted-spvec* (*minus-spmat* *A*)  
 <proof>

**lemma** *sorted-spvec-abs-spmat*: *sorted-spvec* *A*  $\implies$  *sorted-spvec* (*abs-spmat* *A*)  
 <proof>

**lemma** *sorted-spmat-minus-spmat*: *sorted-spmat* *A*  $\implies$  *sorted-spmat* (*minus-spmat* *A*)  
 <proof>

**lemma** *sorted-spmat-abs-spmat*: *sorted-spmat* *A*  $\implies$  *sorted-spmat* (*abs-spmat* *A*)  
 <proof>

**constdefs**

*diff-spmat* :: ('a::lordered-ring) *spmat*  $\Rightarrow$  'a *spmat*  $\Rightarrow$  'a *spmat*



*diff-spmat*  $A\ B == \text{add-spmat } (A, \text{minus-spmat } B)$

**lemma** *sorted-spmat-diff-spmat*: *sorted-spmat*  $A \implies \text{sorted-spmat } B \implies \text{sorted-spmat } (\text{diff-spmat } A\ B)$   
 ⟨proof⟩

**lemma** *sorted-spvec-diff-spmat*: *sorted-spvec*  $A \implies \text{sorted-spvec } B \implies \text{sorted-spvec } (\text{diff-spmat } A\ B)$   
 ⟨proof⟩

**lemma** *sparse-row-diff-spmat*: *sparse-row-matrix*  $(\text{diff-spmat } A\ B) = (\text{sparse-row-matrix } A) - (\text{sparse-row-matrix } B)$   
 ⟨proof⟩

**constdefs**

*sorted-sparse-matrix* :: 'a *spmat*  $\Rightarrow$  *bool*  
*sorted-sparse-matrix*  $A == (\text{sorted-spvec } A) \ \& \ (\text{sorted-spmat } A)$

**lemma** *sorted-sparse-matrix-imp-spvec*: *sorted-sparse-matrix*  $A \implies \text{sorted-spvec } A$   
 ⟨proof⟩

**lemma** *sorted-sparse-matrix-imp-spmat*: *sorted-sparse-matrix*  $A \implies \text{sorted-spmat } A$   
 ⟨proof⟩

**lemmas** *sorted-sp-simps* =  
*sorted-spvec.simps*  
*sorted-spmat.simps*  
*sorted-sparse-matrix-def*

**lemma** *bool1*:  $(\neg \text{True}) = \text{False}$  ⟨proof⟩

**lemma** *bool2*:  $(\neg \text{False}) = \text{True}$  ⟨proof⟩

**lemma** *bool3*:  $((P::\text{bool}) \wedge \text{True}) = P$  ⟨proof⟩

**lemma** *bool4*:  $(\text{True} \wedge (P::\text{bool})) = P$  ⟨proof⟩

**lemma** *bool5*:  $((P::\text{bool}) \wedge \text{False}) = \text{False}$  ⟨proof⟩

**lemma** *bool6*:  $(\text{False} \wedge (P::\text{bool})) = \text{False}$  ⟨proof⟩

**lemma** *bool7*:  $((P::\text{bool}) \vee \text{True}) = \text{True}$  ⟨proof⟩

**lemma** *bool8*:  $(\text{True} \vee (P::\text{bool})) = \text{True}$  ⟨proof⟩

**lemma** *bool9*:  $((P::\text{bool}) \vee \text{False}) = P$  ⟨proof⟩

**lemma** *bool10*:  $(\text{False} \vee (P::\text{bool})) = P$  ⟨proof⟩

**lemmas** *boolarith* = *bool1 bool2 bool3 bool4 bool5 bool6 bool7 bool8 bool9 bool10*

**lemma** *if-case-eq*:  $(\text{if } b \text{ then } x \text{ else } y) = (\text{case } b \text{ of } \text{True} \Rightarrow x \mid \text{False} \Rightarrow y)$   
 ⟨proof⟩

**consts**

*pprt-spvec* :: ('a::{\iordered-ab-group-add}) *spvec*  $\Rightarrow$  'a *spvec*  
*nprt-spvec* :: ('a::{\iordered-ab-group-add}) *spvec*  $\Rightarrow$  'a *spvec*  
*pprt-spmat* :: ('a::{\iordered-ab-group-add}) *spmat*  $\Rightarrow$  'a *spmat*

$np\text{rt-spmat} :: ('a::\{\text{ordered-ab-group-add}\}) \text{ spmat} \Rightarrow 'a \text{ spmat}$

**primrec**

$pp\text{rt-spvec } [] = []$   
 $pp\text{rt-spvec } (a\#as) = (\text{fst } a, pp\text{rt } (\text{snd } a)) \# (pp\text{rt-spvec } as)$

**primrec**

$np\text{rt-spvec } [] = []$   
 $np\text{rt-spvec } (a\#as) = (\text{fst } a, np\text{rt } (\text{snd } a)) \# (np\text{rt-spvec } as)$

**primrec**

$pp\text{rt-spmat } [] = []$   
 $pp\text{rt-spmat } (a\#as) = (\text{fst } a, pp\text{rt-spvec } (\text{snd } a)) \# (pp\text{rt-spmat } as)$

**primrec**

$np\text{rt-spmat } [] = []$   
 $np\text{rt-spmat } (a\#as) = (\text{fst } a, np\text{rt-spvec } (\text{snd } a)) \# (np\text{rt-spmat } as)$

**lemma**  $pp\text{rt-add}$ :  $\text{disj-matrices } A \ (B::(-::\text{ordered-ring}) \text{ matrix}) \Longrightarrow pp\text{rt } (A+B)$   
 $= pp\text{rt } A + pp\text{rt } B$   
 $\langle \text{proof} \rangle$

**lemma**  $np\text{rt-add}$ :  $\text{disj-matrices } A \ (B::(-::\text{ordered-ring}) \text{ matrix}) \Longrightarrow np\text{rt } (A+B)$   
 $= np\text{rt } A + np\text{rt } B$   
 $\langle \text{proof} \rangle$

**lemma**  $pp\text{rt-singleton[simp]}$ :  $pp\text{rt } (\text{singleton-matrix } j \ i \ (x::(-::\text{ordered-ring}))) = \text{singleton-matrix } j \ i \ (pp\text{rt } x)$   
 $\langle \text{proof} \rangle$

**lemma**  $np\text{rt-singleton[simp]}$ :  $np\text{rt } (\text{singleton-matrix } j \ i \ (x::(-::\text{ordered-ring}))) = \text{singleton-matrix } j \ i \ (np\text{rt } x)$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{less-imp-le}$ :  $a < b \Longrightarrow a \leq (b::(-::\text{order})) \langle \text{proof} \rangle$

**lemma**  $\text{sparse-row-vector-pprt}$ :  $\text{sorted-spvec } v \Longrightarrow \text{sparse-row-vector } (pp\text{rt-spvec } v) = pp\text{rt } (\text{sparse-row-vector } v)$   
 $\langle \text{proof} \rangle$

**lemma**  $\text{sparse-row-vector-nprt}$ :  $\text{sorted-spvec } v \Longrightarrow \text{sparse-row-vector } (np\text{rt-spvec } v) = np\text{rt } (\text{sparse-row-vector } v)$   
 $\langle \text{proof} \rangle$

**lemma**  $pp\text{rt-move-matrix}$ :  $pp\text{rt } (\text{move-matrix } (A::('a::\text{ordered-ring}) \text{ matrix}) \ j \ i)$

$= \text{move-matrix } (\text{pprt } A) \ j \ i$   
 $\langle \text{proof} \rangle$

**lemma** *nprrt-move-matrix*:  $\text{nprrt } (\text{move-matrix } (A::('a::\text{lordered-ring}) \text{ matrix}) \ j \ i)$   
 $= \text{move-matrix } (\text{nprrt } A) \ j \ i$   
 $\langle \text{proof} \rangle$

**lemma** *sparse-row-matrix-pprt*:  $\text{sorted-spmat } m \implies \text{sorted-spmat } m \implies \text{sparse-row-matrix}$   
 $(\text{pprt-spmat } m) = \text{pprt } (\text{sparse-row-matrix } m)$   
 $\langle \text{proof} \rangle$

**lemma** *sparse-row-matrix-nprrt*:  $\text{sorted-spmat } m \implies \text{sorted-spmat } m \implies \text{sparse-row-matrix}$   
 $(\text{nprrt-spmat } m) = \text{nprrt } (\text{sparse-row-matrix } m)$   
 $\langle \text{proof} \rangle$

**lemma** *sorted-pprt-spmat*:  $\text{sorted-spmat } v \implies \text{sorted-spmat } (\text{pprt-spmat } v)$   
 $\langle \text{proof} \rangle$

**lemma** *sorted-nprrt-spmat*:  $\text{sorted-spmat } v \implies \text{sorted-spmat } (\text{nprrt-spmat } v)$   
 $\langle \text{proof} \rangle$

**lemma** *sorted-spmat-pprt-spmat*:  $\text{sorted-spmat } m \implies \text{sorted-spmat } (\text{pprt-spmat } m)$   
 $\langle \text{proof} \rangle$

**lemma** *sorted-spmat-nprrt-spmat*:  $\text{sorted-spmat } m \implies \text{sorted-spmat } (\text{nprrt-spmat } m)$   
 $\langle \text{proof} \rangle$

**lemma** *sorted-spmat-pprt-spmat*:  $\text{sorted-spmat } m \implies \text{sorted-spmat } (\text{pprt-spmat } m)$   
 $\langle \text{proof} \rangle$

**lemma** *sorted-spmat-nprrt-spmat*:  $\text{sorted-spmat } m \implies \text{sorted-spmat } (\text{nprrt-spmat } m)$   
 $\langle \text{proof} \rangle$

**constdefs**

$\text{mult-est-spmat} :: ('a::\text{lordered-ring}) \text{ spat} \Rightarrow 'a \text{ spat} \Rightarrow 'a \text{ spat} \Rightarrow 'a \text{ spat}$   
 $\Rightarrow 'a \text{ spat}$   
 $\text{mult-est-spmat } r1 \ r2 \ s1 \ s2 ==$   
 $\text{add-spmat } (\text{mult-spmat } (\text{pprt-spmat } s2) (\text{pprt-spmat } r2), \text{add-spmat } (\text{mult-spmat}$   
 $(\text{pprt-spmat } s1) (\text{nprrt-spmat } r2),$   
 $\text{add-spmat } (\text{mult-spmat } (\text{nprrt-spmat } s2) (\text{pprt-spmat } r1), \text{mult-spmat } (\text{nprrt-spmat}$   
 $s1) (\text{nprrt-spmat } r1))))$

**lemmas** *sparse-row-matrix-op-simps* =

*sorted-sparse-matrix-imp-spmat sorted-sparse-matrix-imp-spmat*  
*sparse-row-add-spmat sorted-spmat-add-spmat sorted-spmat-add-spmat*  
*sparse-row-diff-spmat sorted-spmat-diff-spmat sorted-spmat-diff-spmat*  
*sparse-row-matrix-minus sorted-spmat-minus-spmat sorted-spmat-minus-spmat*

*sparse-row-mult-spmat sorted-spvec-mult-spmat sorted-spmat-mult-spmat*  
*sparse-row-matrix-abs sorted-spvec-abs-spmat sorted-spmat-abs-spmat*  
*le-spmat-iff-sparse-row-le*  
*sparse-row-matrix-pprt sorted-spvec-pprt-spmat sorted-spmat-pprt-spmat*  
*sparse-row-matrix-nprt sorted-spvec-nprt-spmat sorted-spmat-nprt-spmat*

**lemma** *zero-eq-Numeral0*:  $(0::\text{number-ring}) = \text{Numeral0}$  *<proof>*

**lemmas** *sparse-row-matrix-arith-simps*[*simplified zero-eq-Numeral0*] =  
*mult-spmat.simps mult-spvec-spmat.simps*  
*addmult-spvec.simps*  
*smult-spvec-empty smult-spvec-cons*  
*add-spmat.simps add-spvec.simps*  
*minus-spmat.simps minus-spvec.simps*  
*abs-spmat.simps abs-spvec.simps*  
*diff-spmat-def*  
*le-spmat.simps le-spvec.simps*  
*pprt-spmat.simps pprt-spvec.simps*  
*nprrt-spmat.simps nprrt-spvec.simps*  
*mult-est-spmat-def*

**lemma** *spm-mult-le-dual-prts*:

**assumes**

*sorted-sparse-matrix A1*  
*sorted-sparse-matrix A2*  
*sorted-sparse-matrix c1*  
*sorted-sparse-matrix c2*  
*sorted-sparse-matrix y*  
*sorted-sparse-matrix r1*  
*sorted-sparse-matrix r2*  
*sorted-spvec b*  
*le-spmat ([], y)*  
*sparse-row-matrix A1 ≤ A*  
*A ≤ sparse-row-matrix A2*  
*sparse-row-matrix c1 ≤ c*  
*c ≤ sparse-row-matrix c2*  
*sparse-row-matrix r1 ≤ x*  
*x ≤ sparse-row-matrix r2*  
*A \* x ≤ sparse-row-matrix (b::('a::lordered-ring) spmat)*

**shows**

*c \* x ≤ sparse-row-matrix (add-spmat (mult-spmat y b,*  
*(let s1 = diff-spmat c1 (mult-spmat y A2); s2 = diff-spmat c2 (mult-spmat y*  
*A1) in*  
*add-spmat (mult-spmat (pprt-spmat s2) (pprt-spmat r2), add-spmat (mult-spmat*  
*(pprt-spmat s1) (nprrt-spmat r2),*  
*add-spmat (mult-spmat (nprrt-spmat s2) (pprt-spmat r1), mult-spmat (nprrt-spmat*

```

s1) (npmt-spmat r1))))))
⟨proof⟩

```

**lemma** *spm-mult-le-dual-prts-no-let*:

```

assumes
  sorted-sparse-matrix A1
  sorted-sparse-matrix A2
  sorted-sparse-matrix c1
  sorted-sparse-matrix c2
  sorted-sparse-matrix y
  sorted-sparse-matrix r1
  sorted-sparse-matrix r2
  sorted-spmat b
  le-spmat ([], y)
  sparse-row-matrix A1 ≤ A
  A ≤ sparse-row-matrix A2
  sparse-row-matrix c1 ≤ c
  c ≤ sparse-row-matrix c2
  sparse-row-matrix r1 ≤ x
  x ≤ sparse-row-matrix r2
  A * x ≤ sparse-row-matrix (b::('a::lordered-ring) spmat)
shows
  c * x ≤ sparse-row-matrix (add-spmat (mult-spmat y b,
    mult-est-spmat r1 r2 (diff-spmat c1 (mult-spmat y A2)) (diff-spmat c2 (mult-spmat
    y A1))))
⟨proof⟩

```

**end**

**theory** *FloatSparseMatrix* **imports** *Float SparseMatrix* **begin**

**end**

```

theory Compute-Oracle imports Pure
uses am.ML am-compiler.ML am-interpreter.ML am-ghc.ML am-sml.ML report.ML
compute.ML linker.ML
begin

```

⟨*ML*⟩

**end**

```

theory ComputeHOL
imports Main ~~/src/Tools/Compute-Oracle/Compute-Oracle
begin

```

**lemma** *Trueprop-eq-eq*:  $\text{Trueprop } X == (X == \text{True})$   $\langle \text{proof} \rangle$   
**lemma** *meta-eq-trivial*:  $x == y \implies x == y$   $\langle \text{proof} \rangle$   
**lemma** *meta-eq-imp-eq*:  $x == y \implies x = y$   $\langle \text{proof} \rangle$   
**lemma** *eq-trivial*:  $x = y \implies x = y$   $\langle \text{proof} \rangle$   
**lemma** *bool-to-true*:  $x :: \text{bool} \implies x == \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *transmeta-1*:  $x = y \implies y == z \implies x = z$   $\langle \text{proof} \rangle$   
**lemma** *transmeta-2*:  $x == y \implies y = z \implies x = z$   $\langle \text{proof} \rangle$   
**lemma** *transmeta-3*:  $x == y \implies y == z \implies x = z$   $\langle \text{proof} \rangle$

**lemma** *If-True*:  $\text{If } \text{True} = (\lambda x y. x)$   $\langle \text{proof} \rangle$   
**lemma** *If-False*:  $\text{If } \text{False} = (\lambda x y. y)$   $\langle \text{proof} \rangle$

**lemmas** *compute-if* = *If-True If-False*

**lemma** *bool1*:  $(\neg \text{True}) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *bool2*:  $(\neg \text{False}) = \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *bool3*:  $(P \wedge \text{True}) = P$   $\langle \text{proof} \rangle$   
**lemma** *bool4*:  $(\text{True} \wedge P) = P$   $\langle \text{proof} \rangle$   
**lemma** *bool5*:  $(P \wedge \text{False}) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *bool6*:  $(\text{False} \wedge P) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *bool7*:  $(P \vee \text{True}) = \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *bool8*:  $(\text{True} \vee P) = \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *bool9*:  $(P \vee \text{False}) = P$   $\langle \text{proof} \rangle$   
**lemma** *bool10*:  $(\text{False} \vee P) = P$   $\langle \text{proof} \rangle$   
**lemma** *bool11*:  $(\text{True} \longrightarrow P) = P$   $\langle \text{proof} \rangle$   
**lemma** *bool12*:  $(P \longrightarrow \text{True}) = \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *bool13*:  $(\text{True} \longrightarrow P) = P$   $\langle \text{proof} \rangle$   
**lemma** *bool14*:  $(P \longrightarrow \text{False}) = (\neg P)$   $\langle \text{proof} \rangle$   
**lemma** *bool15*:  $(\text{False} \longrightarrow P) = \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *bool16*:  $(\text{False} = \text{False}) = \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *bool17*:  $(\text{True} = \text{True}) = \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *bool18*:  $(\text{False} = \text{True}) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *bool19*:  $(\text{True} = \text{False}) = \text{False}$   $\langle \text{proof} \rangle$

**lemmas** *compute-bool* = *bool1 bool2 bool3 bool4 bool5 bool6 bool7 bool8 bool9 bool10 bool11 bool12 bool13 bool14 bool15 bool16 bool17 bool18 bool19*

**lemma** *compute-fst*:  $\text{fst } (x, y) = x$   $\langle \text{proof} \rangle$   
**lemma** *compute-snd*:  $\text{snd } (x, y) = y$   $\langle \text{proof} \rangle$   
**lemma** *compute-pair-eq*:  $((a, b) = (c, d)) = (a = c \wedge b = d)$   $\langle \text{proof} \rangle$

**lemma** *prod-case-simp*:  $\text{prod-case } f \ (x,y) = f \ x \ y \ \langle \text{proof} \rangle$

**lemmas** *compute-pair* = *compute-fst compute-snd compute-pair-eq prod-case-simp*

**lemma** *compute-the*:  $\text{the } (\text{Some } x) = x \ \langle \text{proof} \rangle$

**lemma** *compute-None-Some-eq*:  $(\text{None} = \text{Some } x) = \text{False} \ \langle \text{proof} \rangle$

**lemma** *compute-Some-None-eq*:  $(\text{Some } x = \text{None}) = \text{False} \ \langle \text{proof} \rangle$

**lemma** *compute-None-None-eq*:  $(\text{None} = \text{None}) = \text{True} \ \langle \text{proof} \rangle$

**lemma** *compute-Some-Some-eq*:  $(\text{Some } x = \text{Some } y) = (x = y) \ \langle \text{proof} \rangle$

**definition**

*option-case-compute* ::  $'b \ \text{option} \Rightarrow 'a \Rightarrow ('b \Rightarrow 'a) \Rightarrow 'a$

**where**

*option-case-compute* *opt a f* = *option-case a f opt*

**lemma** *option-case-compute*:  $\text{option-case} = (\lambda \ a \ f \ \text{opt}. \ \text{option-case-compute } \text{opt } a \ f) \ \langle \text{proof} \rangle$

**lemma** *option-case-compute-None*:  $\text{option-case-compute } \text{None} = (\lambda \ a \ f. \ a) \ \langle \text{proof} \rangle$

**lemma** *option-case-compute-Some*:  $\text{option-case-compute } (\text{Some } x) = (\lambda \ a \ f. \ f \ x) \ \langle \text{proof} \rangle$

**lemmas** *compute-option-case* = *option-case-compute option-case-compute-None option-case-compute-Some*

**lemmas** *compute-option* = *compute-the compute-None-Some-eq compute-Some-None-eq compute-None-None-eq compute-Some-Some-eq compute-option-case*

**lemma** *length-cons*:  $\text{length } (x \# xs) = 1 + (\text{length } xs) \ \langle \text{proof} \rangle$

**lemma** *length-nil*:  $\text{length } [] = 0 \ \langle \text{proof} \rangle$

**lemmas** *compute-list-length* = *length-nil length-cons*

**definition**

*list-case-compute* ::  $'b \ \text{list} \Rightarrow 'a \Rightarrow ('b \Rightarrow 'b \ \text{list} \Rightarrow 'a) \Rightarrow 'a$

**where**

*list-case-compute* *l a f* = *list-case a f l*

**lemma** *list-case-compute*: *list-case* =  $(\lambda (a::'a) f (l::'b \text{ list}). \text{list-case-compute } l \ a \ f)$   
 <proof>

**lemma** *list-case-compute-empty*: *list-case-compute* ( $[]::'b \text{ list}$ ) =  $(\lambda (a::'a) f. a)$   
 <proof>

**lemma** *list-case-compute-cons*: *list-case-compute* ( $u\#v$ ) =  $(\lambda (a::'a) f. (f (u::'b) \ v))$   
 <proof>

**lemmas** *compute-list-case* = *list-case-compute list-case-compute-empty list-case-compute-cons*

**lemma** *compute-list-nth*:  $((x\#xs) ! n) = (\text{if } n = 0 \text{ then } x \text{ else } (xs ! (n - 1)))$   
 <proof>

**lemmas** *compute-list* = *compute-list-case compute-list-length compute-list-nth*

**lemmas** *compute-let* = *Let-def*

**lemmas** *compute-hol* = *compute-if compute-bool compute-pair compute-option compute-list compute-let*

<ML>

**end**

**theory** *ComputeNumeral*  
**imports** *ComputeHOL*  $\sim\sim$  /src/HOL/Real/Float  
**begin**

**lemmas** *bitnorm* = *normalize-bin-simps*

**lemma** *neg1*: *neg Int.Pl* = *False* <proof>  
**lemma** *neg2*: *neg Int.Min* = *True* <proof>  
**lemma** *neg3*: *neg (Int.Bit0 x)* = *neg x* <proof>



**lemma** *neg4*:  $\text{neg } (\text{Int.Bit1 } x) = \text{neg } x$   $\langle \text{proof} \rangle$   
**lemmas** *bitneg* = *neg1 neg2 neg3 neg4*

**lemma** *iszero1*:  $\text{iszero Int.Pls} = \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *iszero2*:  $\text{iszero Int.Min} = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *iszero3*:  $\text{iszero } (\text{Int.Bit0 } x) = \text{iszero } x$   $\langle \text{proof} \rangle$   
**lemma** *iszero4*:  $\text{iszero } (\text{Int.Bit1 } x) = \text{False}$   $\langle \text{proof} \rangle$   
**lemmas** *bitiszero* = *iszero1 iszero2 iszero3 iszero4*

**constdefs**

*lezero*  $x == (x \leq 0)$   
**lemma** *lezero1*:  $\text{lezero Int.Pls} = \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *lezero2*:  $\text{lezero Int.Min} = \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *lezero3*:  $\text{lezero } (\text{Int.Bit0 } x) = \text{lezero } x$   $\langle \text{proof} \rangle$   
**lemma** *lezero4*:  $\text{lezero } (\text{Int.Bit1 } x) = \text{neg } x$   $\langle \text{proof} \rangle$   
**lemmas** *bitlezero* = *lezero1 lezero2 lezero3 lezero4*

**lemma** *biteq1*:  $(\text{Int.Pls} = \text{Int.Pls}) = \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *biteq2*:  $(\text{Int.Min} = \text{Int.Min}) = \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *biteq3*:  $(\text{Int.Pls} = \text{Int.Min}) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *biteq4*:  $(\text{Int.Min} = \text{Int.Pls}) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *biteq5*:  $(\text{Int.Bit0 } x = \text{Int.Bit0 } y) = (x = y)$   $\langle \text{proof} \rangle$   
**lemma** *biteq6*:  $(\text{Int.Bit1 } x = \text{Int.Bit1 } y) = (x = y)$   $\langle \text{proof} \rangle$   
**lemma** *biteq7*:  $(\text{Int.Bit0 } x = \text{Int.Bit1 } y) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *biteq8*:  $(\text{Int.Bit1 } x = \text{Int.Bit0 } y) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *biteq9*:  $(\text{Int.Pls} = \text{Int.Bit0 } x) = (\text{Int.Pls} = x)$   $\langle \text{proof} \rangle$   
**lemma** *biteq10*:  $(\text{Int.Pls} = \text{Int.Bit1 } x) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *biteq11*:  $(\text{Int.Min} = \text{Int.Bit0 } x) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *biteq12*:  $(\text{Int.Min} = \text{Int.Bit1 } x) = (\text{Int.Min} = x)$   $\langle \text{proof} \rangle$   
**lemma** *biteq13*:  $(\text{Int.Bit0 } x = \text{Int.Pls}) = (x = \text{Int.Pls})$   $\langle \text{proof} \rangle$   
**lemma** *biteq14*:  $(\text{Int.Bit1 } x = \text{Int.Pls}) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *biteq15*:  $(\text{Int.Bit0 } x = \text{Int.Min}) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *biteq16*:  $(\text{Int.Bit1 } x = \text{Int.Min}) = (x = \text{Int.Min})$   $\langle \text{proof} \rangle$   
**lemmas** *biteq* = *biteq1 biteq2 biteq3 biteq4 biteq5 biteq6 biteq7 biteq8 biteq9 biteq10 biteq11 biteq12 biteq13 biteq14 biteq15 biteq16*

**lemma** *bitless1*:  $(\text{Int.Pls} < \text{Int.Min}) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *bitless2*:  $(\text{Int.Pls} < \text{Int.Pls}) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *bitless3*:  $(\text{Int.Min} < \text{Int.Pls}) = \text{True}$   $\langle \text{proof} \rangle$   
**lemma** *bitless4*:  $(\text{Int.Min} < \text{Int.Min}) = \text{False}$   $\langle \text{proof} \rangle$   
**lemma** *bitless5*:  $(\text{Int.Bit0 } x < \text{Int.Bit0 } y) = (x < y)$   $\langle \text{proof} \rangle$   
**lemma** *bitless6*:  $(\text{Int.Bit1 } x < \text{Int.Bit1 } y) = (x < y)$   $\langle \text{proof} \rangle$   
**lemma** *bitless7*:  $(\text{Int.Bit0 } x < \text{Int.Bit1 } y) = (x \leq y)$   $\langle \text{proof} \rangle$   
**lemma** *bitless8*:  $(\text{Int.Bit1 } x < \text{Int.Bit0 } y) = (x < y)$   $\langle \text{proof} \rangle$   
**lemma** *bitless9*:  $(\text{Int.Pls} < \text{Int.Bit0 } x) = (\text{Int.Pls} < x)$   $\langle \text{proof} \rangle$

**lemma** *bitless10*:  $(Int.Pls < Int.Bit1\ x) = (Int.Pls \leq x)$  *<proof>*  
**lemma** *bitless11*:  $(Int.Min < Int.Bit0\ x) = (Int.Pls \leq x)$  *<proof>*  
**lemma** *bitless12*:  $(Int.Min < Int.Bit1\ x) = (Int.Min < x)$  *<proof>*  
**lemma** *bitless13*:  $(Int.Bit0\ x < Int.Pls) = (x < Int.Pls)$  *<proof>*  
**lemma** *bitless14*:  $(Int.Bit1\ x < Int.Pls) = (x < Int.Pls)$  *<proof>*  
**lemma** *bitless15*:  $(Int.Bit0\ x < Int.Min) = (x < Int.Pls)$  *<proof>*  
**lemma** *bitless16*:  $(Int.Bit1\ x < Int.Min) = (x < Int.Min)$  *<proof>*  
**lemmas** *bitless* = *bitless1 bitless2 bitless3 bitless4 bitless5 bitless6 bitless7 bitless8*  
*bitless9 bitless10 bitless11 bitless12 bitless13 bitless14 bitless15 bitless16*

**lemma** *bitle1*:  $(Int.Pls \leq Int.Min) = False$  *<proof>*  
**lemma** *bitle2*:  $(Int.Pls \leq Int.Pls) = True$  *<proof>*  
**lemma** *bitle3*:  $(Int.Min \leq Int.Pls) = True$  *<proof>*  
**lemma** *bitle4*:  $(Int.Min \leq Int.Min) = True$  *<proof>*  
**lemma** *bitle5*:  $(Int.Bit0\ x \leq Int.Bit0\ y) = (x \leq y)$  *<proof>*  
**lemma** *bitle6*:  $(Int.Bit1\ x \leq Int.Bit1\ y) = (x \leq y)$  *<proof>*  
**lemma** *bitle7*:  $(Int.Bit0\ x \leq Int.Bit1\ y) = (x \leq y)$  *<proof>*  
**lemma** *bitle8*:  $(Int.Bit1\ x \leq Int.Bit0\ y) = (x < y)$  *<proof>*  
**lemma** *bitle9*:  $(Int.Pls \leq Int.Bit0\ x) = (Int.Pls \leq x)$  *<proof>*  
**lemma** *bitle10*:  $(Int.Pls \leq Int.Bit1\ x) = (Int.Pls \leq x)$  *<proof>*  
**lemma** *bitle11*:  $(Int.Min \leq Int.Bit0\ x) = (Int.Pls \leq x)$  *<proof>*  
**lemma** *bitle12*:  $(Int.Min \leq Int.Bit1\ x) = (Int.Min \leq x)$  *<proof>*  
**lemma** *bitle13*:  $(Int.Bit0\ x \leq Int.Pls) = (x \leq Int.Pls)$  *<proof>*  
**lemma** *bitle14*:  $(Int.Bit1\ x \leq Int.Pls) = (x < Int.Pls)$  *<proof>*  
**lemma** *bitle15*:  $(Int.Bit0\ x \leq Int.Min) = (x < Int.Pls)$  *<proof>*  
**lemma** *bitle16*:  $(Int.Bit1\ x \leq Int.Min) = (x \leq Int.Min)$  *<proof>*  
**lemmas** *bitle* = *bitle1 bitle2 bitle3 bitle4 bitle5 bitle6 bitle7 bitle8*  
*bitle9 bitle10 bitle11 bitle12 bitle13 bitle14 bitle15 bitle16*

**lemmas** *bitsucc* = *succ-bin-simps*

**lemmas** *bitpred* = *pred-bin-simps*

**lemmas** *bituminus* = *minus-bin-simps*

**lemmas** *bitadd* = *add-bin-simps*

**lemma** *mult-Pls-right*:  $x * Int.Pls = Int.Pls$  *<proof>*  
**lemma** *mult-Min-right*:  $x * Int.Min = -\ x$  *<proof>*  
**lemma** *multb0x*:  $(Int.Bit0\ x) * y = Int.Bit0\ (x * y)$  *<proof>*  
**lemma** *multxb0*:  $x * (Int.Bit0\ y) = Int.Bit0\ (x * y)$  *<proof>*  
**lemma** *multb1*:  $(Int.Bit1\ x) * (Int.Bit1\ y) = Int.Bit1\ (Int.Bit0\ (x * y) + x + y)$   
*<proof>*

**lemmas** *bitmul* = *mult-Pls mult-Min mult-Pls-right mult-Min-right multb0x multxb0 multb1*

**lemmas** *bitarith* = *bitnorm bitiszero bitneg bitlezero biteq bitless bitle bitsucc bitpred bituminus bitadd bitmul*

**constdefs**

*nat-norm-number-of* (*x::nat*) == *x*

**lemma** *nat-norm-number-of*: *nat-norm-number-of* (*number-of w*) = (if *lezero w* then 0 else *number-of w*)  
 ⟨*proof*⟩

**lemma** *natnorm0*: (*0::nat*) = *number-of* (*Int.Pls*) ⟨*proof*⟩

**lemma** *natnorm1*: (*1 :: nat*) = *number-of* (*Int.Bit1 Int.Pls*) ⟨*proof*⟩

**lemmas** *natnorm* = *natnorm0 natnorm1 nat-norm-number-of*

**lemma** *natsuc*: *Suc* (*number-of x*) = (if *neg x* then 1 else *number-of* (*Int.succ x*))  
 ⟨*proof*⟩

**lemma** *natadd*: *number-of x* + ((*number-of y*)::*nat*) = (if *neg x* then (*number-of y*) else (if *neg y* then *number-of x* else (*number-of* (*x* + *y*))))  
 ⟨*proof*⟩

**lemma** *natsub*: (*number-of x*) − ((*number-of y*)::*nat*) = (if *neg x* then 0 else (if *neg y* then *number-of x* else (*nat-norm-number-of* (*number-of* (*x* + (− *y*))))))  
 ⟨*proof*⟩

**lemma** *natmul*: (*number-of x*) \* ((*number-of y*)::*nat*) = (if *neg x* then 0 else (if *neg y* then 0 else *number-of* (*x* \* *y*)))  
 ⟨*proof*⟩

**lemma** *nateq*: (((*number-of x*)::*nat*) = (*number-of y*)) = ((*lezero x* ∧ *lezero y*) ∨ (*x* = *y*))  
 ⟨*proof*⟩

**lemma** *natless*: (((*number-of x*)::*nat*) < (*number-of y*)) = ((*x* < *y*) ∧ (¬ (*lezero y*)))  
 ⟨*proof*⟩

**lemma** *natle*: (((*number-of x*)::*nat*) ≤ (*number-of y*)) = (*y* < *x* → *lezero x*)  
 ⟨*proof*⟩

**fun** *natfac* :: *nat*  $\Rightarrow$  *nat*

**where**

*natfac* *n* = (if *n* = 0 then 1 else *n* \* (*natfac* (*n* - 1)))

**lemmas** *compute-natarith* = *bitarith natnorm natsuc natadd natsub natmul nateq*  
*natless natle natfac.simps*

**lemma** *number-eq*: (((*number-of* *x*)::'a::{*number-ring*, *ordered-idom*}) = (*number-of* *y*)) = (*x* = *y*)  
 ⟨*proof*⟩

**lemma** *number-le*: (((*number-of* *x*)::'a::{*number-ring*, *ordered-idom*}) ≤ (*number-of* *y*)) = (*x* ≤ *y*)  
 ⟨*proof*⟩

**lemma** *number-less*: (((*number-of* *x*)::'a::{*number-ring*, *ordered-idom*}) < (*number-of* *y*)) = (*x* < *y*)  
 ⟨*proof*⟩

**lemma** *number-diff*: ((*number-of* *x*)::'a::{*number-ring*, *ordered-idom*}) - *number-of* *y* = *number-of* (*x* + (- *y*))  
 ⟨*proof*⟩

**lemmas** *number-norm* = *number-of-Pls[symmetric] numeral-1-eq-1[symmetric]*

**lemmas** *compute-numberarith* = *number-of-minus[symmetric] number-of-add[symmetric]*  
*number-diff number-of-mult[symmetric] number-norm number-eq number-le number-less*

**lemma** *compute-real-of-nat-number-of*: *real* ((*number-of* *v*)::*nat*) = (if *neg* *v* then 0 else *number-of* *v*)  
 ⟨*proof*⟩

**lemma** *compute-nat-of-int-number-of*: *nat* ((*number-of* *v*)::*int*) = (*number-of* *v*)  
 ⟨*proof*⟩

**lemmas** *compute-num-conversions* = *compute-real-of-nat-number-of compute-nat-of-int-number-of*  
*real-number-of*

**lemmas** *zpowerarith* = *zpower-number-of-even*  
*zpower-number-of-odd[simplified zero-eq-Numeral0-nring one-eq-Numeral1-nring]*  
*zpower-Pls zpower-Min*

**lemma** *adjust*: *adjust* *b* (*q*, *r*) = (if 0 ≤ *r* - *b* then (2 \* *q* + 1, *r* - *b*) else (2 \* *q*, *r*))  
 ⟨*proof*⟩

**lemma** *negateSnd*: *negateSnd* (*q*, *r*) = (*q*, -*r*)

```

    <proof>

lemma divAlg: divAlg (a, b) = (if  $0 \leq a$  then
    if  $0 \leq b$  then posDivAlg a b
    else if  $a=0$  then (0, 0)
    else negateSnd (negDivAlg ( $-a$ ) ( $-b$ ))
  else
    if  $0 < b$  then negDivAlg a b
    else negateSnd (posDivAlg ( $-a$ ) ( $-b$ )))
  <proof>

lemmas compute-div-mod = div-def mod-def divAlg adjust negateSnd posDivAlg.simps
negDivAlg.simps

lemma even-Pls: even (Int.Pls) = True
  <proof>

lemma even-Min: even (Int.Min) = False
  <proof>

lemma even-B0: even (Int.Bit0 x) = True
  <proof>

lemma even-B1: even (Int.Bit1 x) = False
  <proof>

lemma even-number-of: even ((number-of w)::int) = even w
  <proof>

lemmas compute-even = even-Pls even-Min even-B0 even-B1 even-number-of

lemmas compute-numeral = compute-if compute-let compute-pair compute-bool
    compute-natarith compute-numberarith max-def min-def
    compute-num-conversions zpowerarith compute-div-mod compute-even

end

```

```

theory Cplex
imports FloatSparseMatrix ~~/src/HOL/Tools/ComputeNumeral
uses Cplex-tools.ML CplexMatrixConverter.ML FloatSparseMatrixBuilder.ML fspmlp.ML

```

**begin**

**end**

**theory** *MatrixLP*  
**imports** *Cplex*  
**uses** *matrixlp.ML*  
**begin**  
**end**